

First without overflow; when we get to the end of the function, the RSP points to our input ("crypto") on top of the stack and RBP points to the bottom which has:

- 1) the frame pointer (previous RBP)
- 2) return address (where to resume from the calling function) - below the RBP (in the screenshot, but technically above)

```

RBP 0x7fffffffde70 → 0x7fffffffde80 → 0x4011c0 (__libc_csu_init) ← push r15
RSP 0x7fffffffde60 ← 0x6f7470797263 /* 'crypto' */
*RIP 0x40119a (register_name+69) ← leave

[ DISASM ]
0x401185 <register_name+48> mov rsi, rax
0x401188 <register_name+51> lea rdi, [rip+0xebc]
0x40118f <register_name+58> mov eax, 0
0x401194 <register_name+63> call printf@plt
0x401199 <register_name+68> nop
0x40119a <register_name+69> leave
0x40119b <register_name+70> ret
↓
0x4011aa <main+14> mov eax, 0
0x4011af <main+19> pop rbp
0x4011b0 <main+20> ret
↓
0x7ffff7e02d0a <__libc_start_main+234> mov edi, eax

[ STACK ]
00:0000 r10 rsp 0x7fffffffde60 ← 0x6f7470797263 /* 'crypto' */
01:0008 0x7fffffffde68 → 0x401060 (_start) ← xor ebp, ebp
02:0010 rbp 0x7fffffffde70 → 0x7fffffffde80 → 0x4011c0 (__libc_csu_init) ← push 1
03:0018 0x7fffffffde78 → 0x4011aa (main+14) ← mov eax, 0 2
04:0020 0x7fffffffde80 → 0x4011c0 (__libc_csu_init) ← push r15
05:0028 0x7fffffffde88 → 0x7ffff7e02d0a (__libc_start_main+234) ← mov edi,
06:0030 0x7fffffffde90 → 0x7ffff7df78 → 0x7ffff7e29a ← '/home/crystal/Des

```

We hit next and the leave instruction (reverse of "enter" that occurs at beginning of function) resets the stack like:

```

LEAVE
=====
mov %rbp, %rsp    # Sets RBP=RSP, both to same value (0x4011c0 libc_csu_init)
pop %rbp          # Pops the original RBP, that was saved to the stack at start of function

```

Note, because this happens in a single instruction we never actually see RBP = RSP. By the time we get to the next instruction, the value from the top of the stack has been popped into the RBP.

Here we can see that (1) from the last screenshot (the original RBP) has been popped back to the RBP and the address to return to in the calling function is in the RSP (2)

```

*RBP 0x7fffffffde80 → 0x4011c0 (__libc_csu_init) ← push r15 1
*RSP 0x7fffffffde78 → 0x4011aa (main+14) ← mov eax, 0 2
*RIP 0x40119b (register_name+70) ← ret

[ DISASM ]
0x401188 <register_name+51> lea rdi, [rip + 0xebc]
0x40118f <register_name+58> mov eax, 0
0x401194 <register_name+63> call printf@plt <printf>
0x401199 <register_name+68> nop
0x40119a <register_name+69> leave
► 0x40119b <register_name+70> ret <0x40119b>
↓
0x4011aa <main+14> mov eax, 0
0x4011af <main+19> pop rbp
0x4011b0 <main+20> ret
↓
0x7ffff7e02d0a <__libc_start_main+234> mov edi, eax
0x7ffff7e02d0c <__libc_start_main+236> call exit <exit>

[ STACK ]
00:0000 | rsp 0x7fffffffde78 → 0x4011aa (main+14) ← mov eax, 0
01:0008 | rbp 0x7fffffffde80 → 0x4011c0 (__libc_csu_init) ← push r15
02:0010 | 0x7fffffffde88 → 0x7ffff7e02d0a (__libc_start_main+234) ← mov edi, eax
03:0018 | 0x7fffffffde90 → 0x7ffff7ffdf78 → 0x7ffff7ffe29a ← '/home/crystal/Desktop/

```

Now everything is in place to safely return to where we left off in main(). The "ret" instruction will pop the next address (0x4011aa) off the stack into RIP and execute:

```

0x40119a <register_name+69> leave
► 0x40119b <register_name+70> ret
↓
0x4011aa <main+14> mov eax, 0
0x4011af <main+19> pop rbp

```

Now if we repeat for buffer overflow:

```

RBP 0x7fffffffde70 ← 'eaaafaaagaaahaaaiaaajaaakaaalaaamaaanaaaooaaapaaa'
RSP 0x7fffffffde60 ← 'aaaabaaacaaadaaaeaaafaaagaaahaaaiaaajaaakaaalaaamaaanaaaooaaapaaa'
*RIP 0x40119a (register_name+69) ← leave

[ DISASM ]

0x401185 <register_name+48> mov rsi, rax
0x401188 <register_name+51> lea rdi, [rip + 0xebc]
0x40118f <register_name+58> mov eax, 0
0x401194 <register_name+63> call printf@plt <printf@plt>

0x401199 <register_name+68> nop
0x40119a <register_name+69> leave
0x40119b <register_name+70> ret

0x40119c <main> push rbp
0x40119d <main+1> mov rbp, rsp
0x4011a0 <main+4> mov eax, 0
0x4011a5 <main+9> call register_name <register_name>

[ STACK ]

00:0000 | r10 rsp) 0x7fffffffde60 ← 'aaaabaaacaaadaaaeaaafaaagaaahaaaiaaajaaakaaalaaamaaanaaa
01:0008 | 0x7fffffffde68 ← 'caaadaaaeaaafaaagaaahaaaiaaajaaakaaalaaamaaanaaaooaaapaa
02:0010 | {rbp 0x7fffffffde70 ← 'eaaafaaagaaahaaaiaaajaaakaaalaaamaaanaaaooaaapaaa'
03:0018 | 0x7fffffffde78 ← 'gaaahaaaiaaajaaakaaalaaamaaanaaaooaaapaaa'
04:0020 | 0x7fffffffde80 ← 'iaaajaaakaaalaaamaaanaaaooaaapaaa'
05:0028 | 0x7fffffffde88 ← 'kaaalaamaaanaaaooaaapaaa'
06:0030 | 0x7fffffffde90 ← 'maaanaaaooaaapaaa'
07:0038 | 0x7fffffffde98 ← 'oaaapaaa'

```

We've overwritten everything, on the stack including our saved RBP and return pointer from the calling function. So when we LEAVE this time:

LEAVE

=====

```

mov %rbp, %rsp    # Sets RBP=RSP, both to same value (in our case "eaaafaaa...")
pop %rbp          # Pops what *should* be original RBP value (but instead is
"eaaafaaa...") to the RBP

```

Again, we don't see this properly because it happens in one instruction. However, we can see "eaaafaaa" in RBP and the RSP contains the value right after it "gaaahaaa", which would be popped to the RIP and executed (if it were a valid address):

```
*RBP 0x6161616661616165 ('aaaafaaa')
*RSP 0x7fffffffde78 ← 'gaaahaaaiaaaajaaakaaalaaamaaaanaaaooaaapaaa'
*RIP 0x40119b (register_name+70) ← ret

[ DISASM ]
0x401188 <register_name+51> lea rdi, [rip + 0xebc]
0x40118f <register_name+58> mov eax, 0
0x401194 <register_name+63> call printf@plt <printf@plt>
0x401199 <register_name+68> nop
0x40119a <register_name+69> leave
► 0x40119b <register_name+70> ret <0x6161616861616167>

[ STACK ]
00:0000 | rsp 0x7fffffffde78 ← 'gaaahaaaiaaaajaaakaaalaaamaaaanaaaooaaapaaa'
01:0008 | 0x7fffffffde80 ← 'iaaaajaaakaaalaaamaaaanaaaooaaapaaa'
02:0010 | 0x7fffffffde88 ← 'kaaalaaamaaaanaaaooaaapaaa'
03:0018 | 0x7fffffffde90 ← 'maanaaaooaaapaaa'
04:0020 | 0x7fffffffde98 ← 'oaaapaaa'
05:0028 | 0x7fffffffdea0 → 0x401100 (register_tm_clones+48) ← ret
06:0030 | 0x7fffffffdea8 → 0x7ffff7e027cf (init_cacheinfo+287) ← mov rbp, rax
07:0038 | 0x7fffffffdeb0 ← 0x0

[ BACKTRACE ]
```