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(a) Algorithms for Heapsort

Heapsort requires two main algorithms:

Build-Max-Heap: Constructs a max-heap from an unordered array.

Heapify: Maintains the max-heap property.

Heapsort: Sorts the array by repeatedly extracting the maximum element.

Algorithm 1: Build-Max-Heap

Build-Max-Heap(A)

Input: An array A of size n

Output: A max-heap

1. $n = \text{length}(A)$
2. for $i = \lfloor n / 2 \rfloor$ down to 1:
3. Heapify(A, i, n)

Algorithm 2: Heapify

Heapify(A, i, n)

Input: An array A, index i, and heap size n

Output: Maintains max-heap property

1. largest = i
2. left = $2 * i$
3. right = $2 * i + 1$
4. if left $\leq n$ and $A[\text{left}] > A[\text{largest}]$:
5. largest = left
6. if right $\leq n$ and $A[\text{right}] > A[\text{largest}]$:
7. largest = right
8. if largest $\neq i$:
9. Swap $A[i]$ and $A[\text{largest}]$
10. Heapify(A, largest, n)

Algorithm 3: Heapsort

Heapsort(A)

Input: An array A of size n

Output: Sorted array A

1. Build-Max-Heap(A)
 2. $n = \text{length}(A)$
 3. for $i = n$ down to 2:
 4. Swap $A[1]$ and $A[i]$
 5. $n = n - 1$
 6. Heapify(A, 1, n)
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(b) Analysis of Heapsort Algorithms

1. Time Complexity

Heapify: For a node at depth d , Heapify runs in $O(d)O(d)$. Summing over all nodes gives $O(\log n)O(\log n)$.

Build-Max-Heap: The total cost is $O(n)O(n)$, derived from summing $O(\log n)O(\log n)$ for all levels of the heap.

Heapsort: Runs $O(\log n)O(\log n)$ operations n times, resulting in $O(n \log n)O(n \log n)$. Thus, the overall time complexity of Heapsort is $O(n \log n)O(n \log n)$.

2. Space Complexity

Heapsort is in-place and requires no additional space apart from the input array and a few variables. Thus, the space complexity is $O(1)O(1)$.

3. Stability

Heapsort is not stable, as the relative order of equal elements is not preserved during swapping.

(c) Code of Heapsort

```
def heapify(arr, n, i):
    """Maintains the max-heap property."""
    largest = i # Initialize largest as root
    left = 2 * i + 1 # Left child
    right = 2 * i + 2 # Right child

    # Check if left child exists and is greater than root
    if left < n and arr[left] > arr[largest]:
        largest = left

    # Check if right child exists and is greater than largest so far
    if right < n and arr[right] > arr[largest]:
        largest = right

    # If largest is not root, swap and continue heapifying
    if largest != i:
```

```
    arr[i], arr[largest] = arr[largest], arr[i]
    heapify(arr, n, largest)

def build_max_heap(arr):
    """Builds a max-heap from the array."""
    n = len(arr)
    for i in range(n // 2 - 1, -1, -1):
        heapify(arr, n, i)

def heapsort(arr):
    """Sorts the array using Heapsort."""
    n = len(arr)

    # Build a max-heap
    build_max_heap(arr)

    # Extract elements from the heap one by one
    for i in range(n - 1, 0, -1):
        # Move current root to the end
        arr[0], arr[i] = arr[i], arr[0]
        # Call heapify on the reduced heap
        heapify(arr, i, 0)

# Example usage
numbers = [4, 10, 3, 5, 1]
heapsort(numbers)
print("Sorted array:", numbers)
```
