

Optimizaciones

El paper

8.3 The protocol

We describe the protocol below as a non-interactive protocol using the Fiat-Shamir heuristic. For this purpose we always denote by **transcript** the concatenation of the common preprocessed input, and public input, and the proof elements written by the prover up to a certain point in time. We use **transcript** for obtaining random challenges via Fiat-Shamir. One can alternatively, replace all points where we write below “compute challenges”, by the verifier sending random field elements, to obtain the interactive protocol from which we derive the non-interactive one.

Common preprocessed input:

$$\begin{aligned} & n, (x \cdot [1]_1, \dots, x^{n+5} \cdot [1]_1), (q_{Mi}, q_{Li}, q_{Ri}, q_{Oi}, q_{Ci})_{i=1}^n, \sigma^*, \\ & q_M(X) = \sum_{i=1}^n q_{Mi} L_i(X), \\ & q_L(X) = \sum_{i=1}^n q_{Li} L_i(X), \\ & q_R(X) = \sum_{i=1}^n q_{Ri} L_i(X), \\ & q_O(X) = \sum_{i=1}^n q_{Oi} L_i(X), \\ & q_C(X) = \sum_{i=1}^n q_{Ci} L_i(X), \end{aligned}$$

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8.3 The protocol

We describe the protocol below as a non-interactive protocol using the Fiat-Shamir heuristic. For this purpose we always denote by **transcript** the concatenation of the common preprocessed input, and public input, and the proof elements written by the prover up to a certain point in time. We use **transcript** for obtaining random challenges via Fiat-Shamir. One can alternatively, replace all points where we write below “compute challenges”, by the verifier sending random field elements, to obtain the interactive protocol from which we derive the non-interactive one.

Common preprocessed input:

SRS

$$n, (x \cdot [1]_1, \dots, x^{n+5} \cdot [1]_1), (q_{Mi}, q_{Li}, q_{Ri}, q_{Oi}, q_{Ci})_{i=1}^n, \sigma^*,$$

$$q_M(X) = \sum_{i=1}^n q_{Mi} L_i(X),$$

$$q_L(X) = \sum_{i=1}^n q_{Li} L_i(X),$$

$$q_R(X) = \sum_{i=1}^n q_{Ri} L_i(X),$$

$$q_O(X) = \sum_{i=1}^n q_{Oi} L_i(X),$$

$$q_C(X) = \sum_{i=1}^n q_{Ci} L_i(X),$$



El paper

$$\begin{aligned} S_{\sigma 1}(X) &= \sum_{i=1}^n \sigma^*(i) L_i(X), \\ S_{\sigma 2}(X) &= \sum_{i=1}^n \sigma^*(n+i) L_i(X), \\ S_{\sigma 3}(X) &= \sum_{i=1}^n \sigma^*(2n+i) L_i(X) \end{aligned}$$

Compute linearisation polynomial $r(X)$:

$$\begin{aligned} r(X) = & [\bar{a}\bar{b} \cdot q_M(X) + \bar{a} \cdot q_L(X) + \bar{b} \cdot q_R(X) + \bar{c} \cdot q_O(X) + Pl(\mathfrak{z}) + q_C(X)] \\ & + \alpha[(\bar{a} + \beta\mathfrak{z} + \gamma)(\bar{b} + \beta k_1\mathfrak{z} + \gamma)(\bar{c} + \beta k_2\mathfrak{z} + \gamma) \cdot z(X) \\ & - (\bar{a} + \beta\bar{s}_{\sigma 1} + \gamma)(\bar{b} + \beta\bar{s}_{\sigma 2} + \gamma)(\bar{c} + \beta \cdot S_{\sigma 3}(X) + \gamma)\bar{z}_\omega] \\ & + \alpha^2[(z(X) - 1)L_1(\mathfrak{z})] \\ & - Z_H(\mathfrak{z}) \cdot (t_{lo}(X) + \mathfrak{z}^n t_{mid}(X) + \mathfrak{z}^{2n} t_{hi}(X)) \end{aligned}$$

$$a(X) = (b_1 X + b_2) Z_H(X) + \sum_{i=1}^n w_i L_i(X)$$

$$b(X) = (b_3 X + b_4) Z_H(X) + \sum_{i=1}^n w_{n+i} L_i(X)$$

$$c(X) = (b_5 X + b_6) Z_H(X) + \sum_{i=1}^n w_{2n+i} L_i(X)$$

Compute permutation polynomial $z(X)$:

$$\begin{aligned} z(X) = & (b_7 X^2 + b_8 X + b_9) Z_H(X) \\ & + L_1(X) + \sum_{i=1}^{n-1} \left(L_{i+1}(X) \prod_{j=1}^i \frac{(w_j + \beta \omega^j + \gamma)(w_{n+j} + \beta k_1 \omega^j + \gamma)(w_{2n+j} + \beta k_2 \omega^j + \gamma)}{(w_j + \sigma^*(j)\beta + \gamma)(w_{n+j} + \sigma^*(n+j)\beta + \gamma)(w_{2n+j} + \sigma^*(2n+j)\beta + \gamma)} \right) \end{aligned}$$

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$$a(X) = (b_1 X + b_2) Z_H(X) + \sum_{i=1}^n w_i L_i(X)$$



Compute

Compu

$$+L_1(X) + \sum_{i=1}^n \left(L_{i+1}(X) \prod_{j=1}^i \frac{(w_j + \sigma^*(j)\beta + \gamma)(w_{n+j} + \sigma^*(n+j)\beta + \gamma)(w_{2n+j} + \sigma^*(2n+j)\beta + \gamma)}{(w_j + \sigma^*(j)\beta + \gamma)(w_{n+j} + \sigma^*(n+j)\beta + \gamma)(w_{2n+j} + \sigma^*(2n+j)\beta + \gamma)} \right)$$



Entendiendo el paper

Interpolación de la traza

$$a(X) = (b_1X + b_2)Z_H(X) + \sum_{i=1}^n w_i L_i(X)$$

$$b(X) = (b_3X + b_4)Z_H(X) + \sum_{i=1}^n w_{n+i} L_i(X)$$

$$c(X) = (b_5X + b_6)Z_H(X) + \sum_{i=1}^n w_{2n+i} L_i(X)$$

Interpolación de la traza

$$a(X) = \text{[gray box]} + \sum_{i=1}^n w_i L_i(X)$$

- Base de Lagrange:

$$L_i(g^i) = 1$$

$$L_i(g^j) = 0 \text{ para todo } j \text{ distinto de } i$$

- $a(g^i) = w_i$

Blindings (ZK!)

$$Z_H = X^n - 1$$

$$a(X) = (b_1X + b_2)Z_H(X) + \sum_{i=1}^n w_i L_i(X)$$

$$b(X) = (b_3X + b_4)Z_H(X) + \sum_{i=1}^n w_{n+i} L_i(X)$$

$$c(X) = (b_5X + b_6)Z_H(X) + \sum_{i=1}^n w_{2n+i} L_i(X)$$

Blindings (ZK!)

$$Z_H = X^n - 1$$

$$a(X) = (b_1X + b_2)Z_H(X) + \sum_{i=1}^n w_i L_i(X)$$

- No modifica los valores en el dominio $a(g^i) = w_i$
- Randomiza los valores por fuera del dominio.
 - Consultar $a(z_1)$ y $a(z_2)$ no revela info sobre la traza

El S_σ está raro...

$$\begin{aligned} S_{\sigma 1}(X) &= \sum_{i=1}^n \sigma^*(i) L_i(X), \\ S_{\sigma 2}(X) &= \sum_{i=1}^n \sigma^*(n+i) L_i(X), \\ S_{\sigma 3}(X) &= \sum_{i=1}^n \sigma^*(2n+i) L_i(X) \end{aligned}$$



Cuando vimos el protocolo para las máscaras **con la traza aplanada**:

$$Z(X)(v_1(X) + \alpha) - Z(gX)(v_2(X) + \alpha)$$

$$v_1(X) \rightsquigarrow \beta D + W$$

$$v_2(X) \rightsquigarrow \beta \sigma(D) + W$$

El S_σ está raro...

$$\begin{aligned} S_{\sigma 1}(X) &= \sum_{i=1}^n \sigma^*(i) L_i(X), \\ S_{\sigma 2}(X) &= \sum_{i=1}^n \sigma^*(n+i) L_i(X), \\ S_{\sigma 3}(X) &= \sum_{i=1}^n \sigma^*(2n+i) L_i(X) \end{aligned}$$



Cuando vimos el protocolo para las máscaras **con la traza W aplanada**:

$$Z(X)(v_1(X) + \alpha) - Z(gX)(v_2(X) + \alpha)$$

$$v_1(X) = \beta X + w(X)$$

$$v_2(X) = \beta S_\sigma(X) + w(X)$$

El S_σ está raro...

$$\begin{aligned} S_{\sigma 1}(X) &= \sum_{i=1}^n \sigma^*(i) L_i(X), \\ S_{\sigma 2}(X) &= \sum_{i=1}^n \sigma^*(n+i) L_i(X), \\ S_{\sigma 3}(X) &= \sum_{i=1}^n \sigma^*(2n+i) L_i(X) \end{aligned}$$



Cuando vimos el protocolo para las máscaras **con la traza W aplanada**:

$$Z(X)(\beta X + w(X) + \alpha) - Z(gX)(\beta S_\sigma(X) + w(X) + \alpha)$$

El S_σ está raro...

$$\begin{aligned} S_{\sigma 1}(X) &= \sum_{i=1}^n \sigma^*(i) L_i(X), \\ S_{\sigma 2}(X) &= \sum_{i=1}^n \sigma^*(n+i) L_i(X), \\ S_{\sigma 3}(X) &= \sum_{i=1}^n \sigma^*(2n+i) L_i(X) \end{aligned}$$

Sin aplanar esto queda así:

$$\begin{aligned} &Z(X)(\beta X + a(X) + \alpha)(\beta X + b(X) + \alpha)(\beta X + c(X) + \alpha) \\ &\quad - Z(gX)(\beta S_{\sigma 1}(X) + a(X) + \alpha)(\beta S_{\sigma 2}(X) + b(X) + \alpha)(\beta S_{\sigma 3}(X) + c(X) + \alpha) \end{aligned}$$



Juntar los dos protocolos

Vimos protocolos para probar que

- La traza resuelve las ecuaciones Q fila a fila

$$a(X)q_L(X) + b(X)q_R(X) + a(X)b(X)q_M(X) + c(X)q_O(X) + q_C(X)$$

- La traza respeta la máscara

$$\begin{aligned} & Z(X)(\beta X + a(X) + \alpha)(\beta X + b(X) + \alpha)(\beta X + c(X) + \alpha) \\ & - Z(gX)(\beta S_{\sigma_1}(X) + a(X) + \alpha)(\beta S_{\sigma_2}(X) + b(X) + \alpha)(\beta S_{\sigma_3}(X) + c(X) + \alpha) \end{aligned}$$

Juntar los dos protocolos

Un solo polinomio haciendo una combinación lineal aleatoria **con γ random**

$$a(X)q_L(X) + b(X)q_R(X) + a(X)b(X)q_M(X) + c(X)q_O(X) + q_C(X)$$

$$+ \gamma \left(\begin{aligned} &Z(X)(\beta X + a(X) + \alpha)(\beta X + b(X) + \alpha)(\beta X + c(X) + \alpha) \\ &- Z(gX)(\beta S_{\sigma_1}(X) + a(X) + \alpha)(\beta S_{\sigma_2}(X) + b(X) + \alpha)(\beta S_{\sigma_3}(X) + c(X) + \alpha) \end{aligned} \right)$$

Juntar los dos protocolos

después de dividir por $Z_H = X^n - 1$

$$\begin{aligned} t(X) = & (a(X)b(X)q_M(X) + a(X)q_L(X) + b(X)q_R(X) + c(X)q_O(X) + Pl(X) + q_C(X)) \frac{1}{Z_H(X)} \\ & + ((a(X) + \beta X + \gamma)(b(X) + \beta k_1 X + \gamma)(c(X) + \beta k_2 X + \gamma)z(X)) \frac{\alpha}{Z_H(X)} \\ & - ((a(X) + \beta S_{\sigma_1}(X) + \gamma)(b(X) + \beta S_{\sigma_2}(X) + \gamma)(c(X) + \beta S_{\sigma_3}(X) + \gamma)z(X\omega)) \frac{\alpha}{Z_H(X)} \end{aligned}$$



En el paper α es γ y γ es α

Pero en realidad...

$$\begin{aligned} t(X) = & (a(X)b(X)q_M(X) + a(X)q_L(X) + b(X)q_R(X) + c(X)q_O(X) + Pl(X) + q_C(X)) \frac{1}{z_H(X)} \\ & + ((a(X) + \beta X + \gamma)(b(X) + \beta k_1 X + \gamma)(c(X) + \beta k_2 X + \gamma)z(X)) \frac{\alpha}{z_H(X)} \\ & - ((a(X) + \beta S_{\sigma 1}(X) + \gamma)(b(X) + \beta S_{\sigma 2}(X) + \gamma)(c(X) + \beta S_{\sigma 3}(X) + \gamma)z(X\omega)) \frac{\alpha}{z_H(X)} \\ & + (z(X) - 1) L_1(X) \frac{\alpha^2}{z_H(X)} \end{aligned}$$

Pero en realidad...

$$\begin{aligned} t(X) = & (a(X)b(X)q_M(X) + a(X)q_L(X) + b(X)q_R(X) + c(X)q_O(X) + Pl(X) + q_C(X)) \frac{1}{Z_H(X)} \\ & + ((a(X) + \beta X + \gamma)(b(X) + \beta k_1 X + \gamma)(c(X) + \beta k_2 X + \gamma)z(X)) \frac{\alpha}{Z_H(X)} \\ & - ((a(X) + \beta S_{\sigma 1}(X) + \gamma)(b(X) + \beta S_{\sigma 2}(X) + \gamma)(c(X) + \beta S_{\sigma 3}(X) + \gamma)z(X\omega)) \frac{\alpha}{Z_H(X)} \\ & + (z(X) - 1) L_1(X) \frac{\alpha^2}{Z_H(X)} \end{aligned}$$

**Equivalente a chequear Z(1)
distinto de cero sin hacer
openings**

Y eso?



$$\begin{aligned}
 t(X) = & (a(X)b(X)q_M(X) + a(X)q_L(X) + b(X)q_R(X) + c(X)q_O(X) + \boxed{PI(X)} + q_C(X)) \frac{1}{z_H(X)} \\
 & + ((a(X) + \beta X + \gamma)(b(X) + \beta k_1 X + \gamma)(c(X) + \beta k_2 X + \gamma)z(X)) \frac{\alpha}{z_H(X)} \\
 & - ((a(X) + \beta S_{\sigma_1}(X) + \gamma)(b(X) + \beta S_{\sigma_2}(X) + \gamma)(c(X) + \beta S_{\sigma_3}(X) + \gamma)z(X\omega)) \frac{\alpha}{z_H(X)} \\
 & + (z(X) - 1) L_1(X) \frac{\alpha^2}{z_H(X)}
 \end{aligned}$$

Partiendo t

Common preprocessed input:

x es nuestro τ

$$n) (x \cdot [1]_1, \dots, x^{n+5} \cdot [1]_1), (q_{Mi}, q_i)$$

$$q_M(X) = \sum_{i=1}^n q_{Mi} L_i(X),$$

$$a_i(X) = \sum_{j=1}^n a_{ij} L_j(X).$$

Split $t(X)$ into degree $< n$ polynomials $t'_{lo}(X)$, $t'_{mid}(X)$ and $t'_{hi}(X)$ of degree at most $n + 5$, such that

$$t(X) = t'_{lo}(X) + X^n t'_{mid}(X) + X^{2n} t'_{hi}(X)$$

Now choose random scalars $b_{10}, b_{11} \in \mathbb{F}$ and define

$$t_{lo}(X) := t'_{lo}(X) + b_{10} X^n, t_{mid}(X) := t'_{mid}(X) - b_{10} + b_{11} X^n, t_{hi}(X) := t'_{hi}(X) - b_{11}$$

Partiendo t



Ariel Gabizon

@rel_zeta_tech

To all plonkers out there.

A talented student from TU Wien named Marek Sefranek has discovered a mistake in the implementation of zero-knowledge in Section 8 of the plonk paper.

8:44 AM · Jun 30, 2022

Paper 2024/848

How (Not) to Simulate PLONK

Marek Sefranek , TU Wien

Split $t(X)$ into degree $< n$ polynomials $t_{\text{lo}}, t_{\text{mid}}, t_{\text{hi}}$ of degree $< n + 5$, such that

$$t(X) = t'_{\text{lo}}(X) + X^n t'_{\text{mid}}(X) + X^{2n} t'_{\text{hi}}(X)$$

Now choose random scalars $b_{10}, b_{11} \in \mathbb{F}$ and define

$$t_{\text{lo}}(X) := t'_{\text{lo}}(X) + b_{10}X^n, t_{\text{mid}}(X) := t'_{\text{mid}}(X) - b_{10} + b_{11}X^n, t_{\text{hi}}(X) := t'_{\text{hi}}(X) - b_{11}$$

Linearisation trick

Supongamos que ya commiteamos $[h_1], [h_2], [h_3]$

Y queremos probar

$$h_1(X)h_2(X) - h_3(X) = 0$$

Linearisation trick

Supongamos que ya committeamos $[h_1], [h_2], [h_3]$

Y queremos probar

$$h_1(X)h_2(X) - h_3(X) = 0$$

Opción 1: Enviar $h_1(z), h_2(z), h_3(z)$ con sus respectivas pruebas de KZG

Linearisation trick

Supongamos que ya commiteamos $[h_1], [h_2], [h_3]$

Y queremos probar

$$h_1(X)h_2(X) - h_3(X) = 0$$

Opción 1: Enviar $h_1(z)$, $h_2(z)$, $h_3(z)$ con sus respectivas pruebas de KZG

Opción 2:

1. Enviar $h_2(z)$ con su respectiva prueba de KZG.
2. Enviar la prueba de KZG de que $h_1(X)h_2(z) - h_3(X)$ vale 0 en z

Linearisation trick

Compute linearisation polynomial $r(X)$:

$$\begin{aligned} r(X) = & \left[\bar{a}\bar{b} \cdot \mathbf{q}_M(X) + \bar{a} \cdot \mathbf{q}_L(X) + \bar{b} \cdot \mathbf{q}_R(X) + \bar{c} \cdot \mathbf{q}_O(X) + \text{Pl}(\mathfrak{z}) + \mathbf{q}_C(X) \right] \\ & + \alpha [(\bar{a} + \beta \mathfrak{z} + \gamma)(\bar{b} + \beta k_1 \mathfrak{z} + \gamma)(\bar{c} + \beta k_2 \mathfrak{z} + \gamma) \cdot \mathbf{z}(X) \\ & - (\bar{a} + \beta \bar{\mathbf{s}}_{\sigma 1} + \gamma)(\bar{b} + \beta \bar{\mathbf{s}}_{\sigma 2} + \gamma)(\bar{c} + \beta \cdot \mathbf{S}_{\sigma 3}(X) + \gamma) \bar{z}_\omega] \\ & + \alpha^2 [(\mathbf{z}(X) - 1)\mathbf{L}_1(\mathfrak{z})] \\ & - Z_H(\mathfrak{z}) \cdot (t_{lo}(X) + \mathfrak{z}^n t_{mid}(X) + \mathfrak{z}^{2n} t_{hi}(X)) \end{aligned}$$



Batch Opening (1)

Si necesitamos commitear polinomios f_0, \dots, f_k para luego abrirlos en un elemento random ζ

Basta con commitear solo

$$f_0 + v f_1 + \dots + v^k f_k$$

donde v es otro random elegido por el verifier

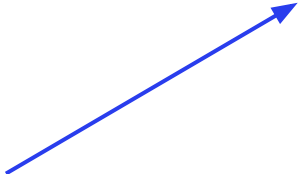
Batch Opening (1)

Compute opening proof polynomial $W_{\mathfrak{z}}(X)$:

$$W_{\mathfrak{z}}(X) = \frac{1}{X - \mathfrak{z}} \begin{pmatrix} r(X) \\ +v(a(X) - \bar{a}) \\ +v^2(b(X) - \bar{b}) \\ +v^3(c(X) - \bar{c}) \\ +v^4(S_{\sigma 1}(X) - \bar{s}_{\sigma 1}) \\ +v^5(S_{\sigma 2}(X) - \bar{s}_{\sigma 2}) \end{pmatrix}$$

Batch Opening (1)

Compute opening proof polynomial $W_{\mathfrak{z}}(X)$:

$$W_{\mathfrak{z}}(X) = \frac{1}{X - \mathfrak{z}} \begin{pmatrix} r(X) \\ +v(a(X) - \bar{a}) \\ +v^2(b(X) - \bar{b}) \\ +v^3(c(X) - \bar{c}) \\ +v^4(S_{\sigma 1}(X) - \bar{s}_{\sigma 1}) \\ +v^5(S_{\sigma 2}(X) - \bar{s}_{\sigma 2}) \end{pmatrix}$$


Polinomio auxiliar en el batch Open de KZG

Batch Opening (2)

Si necesitamos commitear polinomios f_1, f_2 para luego abrirlos en elementos elemento random ζ_1 y ζ_2

Basta con chequear una sola igualdad de pairings

$$e([t_1] + u[t_2], [\tau]) = e(\zeta_1[t_1] + u\zeta_2[t_2] + [f_1 - y_1] + [f_2 - y_2], [1])$$

Batch Opening (2)

$$e([t_1] + u[t_2], [\tau]) = e(\zeta_1[t_1] + u\zeta_2[t_2] + [f_1 - y_1] + [f_2 - y_2], [1])$$

12. Batch validate all evaluations:

$$e([W_3]_1 + u \cdot [W_{3\omega}]_1, [x]_2) \stackrel{?}{=} e(\mathfrak{z} \cdot [W_3]_1 + u\mathfrak{z}\omega \cdot [W_{3\omega}]_1 + [F]_1 - [E]_1, [1]_2)$$

¡Muchas gracias!