

# R1CS Programming

## ZK0x04 Workshop Notes

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## 1 Multiplicative inverse

Deterministically computing  $1/x$  in an R1CS circuit would be expensive. Instead, we can have the prover compute  $1/x$  outside of the circuit and supply the result as a witness element, which we will call  $x_{\text{inv}}$ . To verify the result, we enforce

$$(x)(x_{\text{inv}}) = (1) \quad (1)$$

## 2 Zero testing

To assert  $x = 0$ , we simply enforce

$$(x)(1) = (0) \quad (2)$$

Asserting  $x \neq 0$  is similarly easy: we compute  $1/x$  (non-deterministically, as in [Section 1](#)). The result can be ignored; the mere fact that an inverse exists implies  $x \neq 0$ .

On the other hand, if we want to *evaluate*

$$y = \begin{cases} 0 & \text{if } x = 0 \\ 1 & \text{otherwise} \end{cases} \quad (3)$$

we can do so by introducing another variable  $m$ , and enforcing

$$(x)(m) = (y) \quad (4)$$

$$(1 - y)(x) = (0) \quad (5)$$

TODO: Explain further. This method is from [\[1\]](#).

## 3 Binary

To assert  $b \in \{0, 1\}$ , we enforce

$$(b)(b - 1) = (0) \quad (6)$$

## 4 Comparisons

TODO: Describe basic comparison algorithm

TODO: Describe Ahmed's optimization

A few other optimizations are possible in particular circumstances:

1. To assert (not evaluate)  $x < y$ , we can split  $x$  non-canonically and split  $y$  canonically. The prover is forced to use  $x$ 's canonical representation anyway, otherwise  $x_{\text{bin}} \geq |F| > y_{\text{bin}}$ , making the assertion unsatisfiable.
2. To assert  $x < c$  for some constant  $c \ll |F|$ , we can split  $x$  into just  $\lceil \log_2(c) \rceil$  bits.

## 5 Permutations

Say we want to verify that two sequences,  $(x_1, \dots, x_n)$  and  $(y_1, \dots, y_n)$ , are permutations of one another.

## 6 Sorting

TODO: Discuss sorting networks

TODO: Discuss permutation networks + comparisons to verify order

## 7 Random access

TODO: Discuss naive random access via index comparisons

TODO: Discuss binary tree method

## 8 Embedded curve operations

TODO: Discuss basic embedded curve operations

## References

- [1] B. Parno, J. Howell, C. Gentry, and M. Raykova, “Pinocchio: Nearly practical verifiable computation,” in *2013 IEEE Symposium on Security and Privacy*, pp. 238–252, IEEE, 2013.