R1CS Programming ZK0x04 Workshop Notes

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1 Multiplicative inverse

Deterministically computing 1/x in an R1CS circuit would be expensive. Instead, we can have the prover compute 1/x outside of the circuit and supply the result as a witness element, which we will call $x_{\rm inv}$. To verify the result, we enforce

$$(x)(x_{\rm inv}) = (1) \tag{1}$$

2 Zero testing

To assert x = 0, we simply enforce

$$(x)(1) = (0) (2)$$

Asserting $x \neq 0$ is similarly easy: we compute 1/x (non-deterministically, as in Section 1). The result can be ignored; the mere fact that an inverse exists implies $x \neq 0$.

On the other hand, if we want to evaluate

$$y = \begin{cases} 0 & \text{if } x = 0\\ 1 & \text{otherwise} \end{cases}$$
 (3)

we can do so by introducing another variable m, and enforcing

$$(x)(m) = (y) (4)$$

$$(1-y)(x) = 0 (5)$$

This method is from [1].

3 Binary

To assert $b \in \{0,1\}$, we enforce

$$(b)(b-1) = (0) (6)$$

4 Comparisons

TODO: Describe basic comparison algorithm

TODO: Describe Ahmed's optimization

A few other optimizations are possible in particular circumstances:

- 1. To assert (not evaluate) x < y, we can split x non-canonically and split y canonically. The prover is forced to use x's canonical representation anyway, otherwise $x_{\rm bin} \ge |F| > y_{\rm bin}$, making the assertion unsatisfiable.
- 2. To assert x < c for some constant $c \ll |F|$, we can split x into just $\lceil \log_2(c) \rceil$ bits.

5 Permutations

Say we want to verify that two sequences, (x_1, \ldots, x_n) and (y_1, \ldots, y_n) , are permutations of one another.

6 Sorting

TODO: Discuss sorting networks

TODO: Discuss permutation networks + comparisons to verify order

7 Random access

TODO: Discuss naive random access via index comparisons

TODO: Discuss binary tree method

8 Embedded curve operations

TODO: Discuss basic embedded curve operations

References

[1] B. Parno, J. Howell, C. Gentry, and M. Raykova, "Pinocchio: Nearly practical verifiable computation," in 2013 IEEE Symposium on Security and Privacy, pp. 238–252, IEEE, 2013.