

Shellcode: advantages

- Self-sufficient: easy to inject into other applications
- Small: can fit into a tiny space i.e. caves between sections
- May be used as a loader: first code injected into an application, that follows to load other modules
- Sometimes (but less often) the full malicious functionality can be implemented as shellcode (i.e. Fobber malware)
- This type of code was popular in the past, virus era: where malware code was added to existing PE files (rather than injected into processes)

Creating shellcode

- In case of PE format we just write a code and don't have to worry how it is loaded: Windows Loader will do it
- It is different when we write shellcode
- We cannot rely on the conviniences provided by PE format and Windows Loader:
 - No sections
 - No Data Directories (imports, relocations)
 - Only code to provide everything we need...

Creating shellcode

Feature	PE file	shellcode
Loading	 via Windows Loader running new EXE triggers creation of a new process 	 Custom, simplified must parasite on existing process (i.e. via code injection + thread injection)
Composition	Sections with specific access rights, carrying various elements (code, data, resources, etc)	All in one memory area (read,write,execute)
Relocation to the load base	Defined by relocation table, applied by Windows Loader	Cutom; position-independent code
Access to system API (Imports loading)	Defined by import table, applied by Windows Loader	Custom: retrieving imports via PEB lookup; no IAT, or simplified

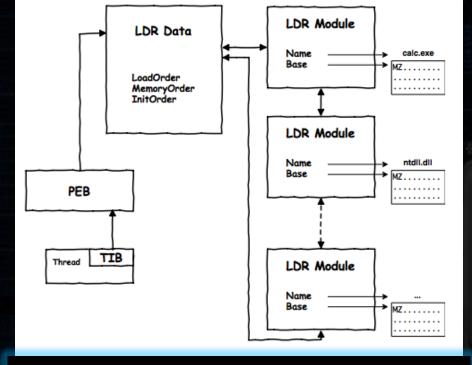
Position-independent code

- In order to create a position-independent code, we must take care that all the addresses that we use are relative to the current intruction pointer address
- A short jump, long jump, call to a local function are relative -> we can use them!

```
EBEO V JMP SHORT 0X413BA8
E8EE000000 V CALL 0X413BFA
```

Any address that needs to be relocated (i.e. using of the data from different PE section)
 breaks the position independence:

• In order to retrieve the imported functions, we will take advantage of the linklist pointed by PEB





http://blog.malcom.pl/2017/shellcode-peb-i-adres-bazowy-modulu-kernel32-dll.html

 In order to retrieve the imported functions, we will take advantage of the linklist pointed by PEB

```
Pyoid Reserved1[12]:
    PPEB ProcessEnvironmentBlock;
    PVOID Reserved2[399];
    BYTE Reserved3[1952];
    PVOID TlsSlots[64];
    BYTE Reserved4[8];
    PVOID Reserved5[26];
    PVOID Reserved5[26];
    PVOID Reserved6[4];
    PVOID TlsExpansionSlots;
} TEB, *PTEB;
```

```
typedef struct PEB
     UCHAR InheritedAddressSpace;
    UCHAR ReadImageFileExecOptions:
     UCHAR BeingDebugged;
    UCHAR BitField:
    ULONG ImageUsesLargePages: 1:
     ULONG IsProtectedProcess: 1:
    ULONG IsLegacyProcess: 1;
    ULONG IsImageDynamicallyRelocated: 1;
    ULONG SpareBits: 4;
     PVOID Mutant:
     PVOID ImageBaseAddress:
    PPEB LDR DATA Ldr;
     FLS CALLBACK INFO * FlsCallback;
    LIST ENTRY FlsListHead;
     PVOID FlsBitmap;
    ULONG FlsBitmapBits[4];
    ULONG FlsHighIndex;
     PVOID WerRegistrationData;
     PVOID WerShipAssertPtr;
 PEB, *PPEB;
```

```
typedef struct _PEB_LDR_DATA

{
    ULONG Length;
    BOOLEAN Initialized;
    HANDLE SSHandle:
    LIST_ENTRY InLoadOrderModuleList;
    LIST_ENTRY InMemoryOrderModuleList;
    LIST_ENTRY InInitializationOrderModuleList;
    PVOID EntryInProgress;
}

PEB_LDR_DATA, *PPEB_LDR_DATA;
```

• We will process each entry, searching for the DLL that we need...

```
typedef struct _PEB_LDR_DATA

{
    ULONG Length;
    BOOLEAN Initialized;
    HANDLE SsHandle;
    LIST_ENTRY InLoadOrderModuleList;
    LIST_ENTRY InMemoryOrderModuleList;
    LIST_ENTRY InInitializationOrderModuleList;
    PVOID EntryInProgress;
}

} PEB_LDR_DATA, *PPEB_LDR_DATA;
```

```
lkd> dt ntdll! LDR DATA TABLE ENTRY
   +0x000 InLoadOrderLinks : LIST ENTRY
   +0x010 InMemoryOrderLinks : _LISI_ENIRY
                                     : LIST ENTRY
   +0x020 InInitializationOrderLinks
   +0x030 DllBase
                           : Ptr64 Void
   +0x038 EntryPoint
                             Ptr64 Void
   +0x040 SizeOfImage
                             Uint4B
  +0x048 FullDilName
                              HNICODE STRING
   +0x058 BaseDllName
                              UNICODE STRING
   +0x068 FlagGroup
                               l] UChar
```

Next LDR_DATA_TABLE_ENTRY

typedef struct _UNICODE_STRING {
 USHORT Length;
 USHORT MaximumLength;
 PWSTR Buffer;
} UNICODE_STRING, *PUNICODE_STRING;

L"Ntdll.dll"

- 1. Get the PEB address
- 2. Via PEB->Ldr->InMemoryOrderModuleList, find:
 - kernel32.d11 (loaded in majority of the processes after initialization)
 - or ntdll·dll (if we want to use low-leven equivalents of Import loading functions)
- 3. Walk through exports table to find addresses of:
 - LoadLibraryA (eventually: ntdll-LdrLoadDll)
 - GetProcAddress (eventually: ntdll LdrGetProcedureAddress)
- 4. Use LoadLibrary A to load other needed DLLs
- 5. Use GetProcAddress to retrieve functions

Creating shellcode: assembly

• We can use YASM for shellcodes written in pure assembly:

```
yasm -f bin demo∙asm
```

- We will not use a linker, which means:
 - we need to fill imports by ourselves
 - we need to take care of relocations or make the code position-independent

Creating shellcode: C

• We can use a C compiler to generate assembly:

```
C1 /c /FA <file_name>.cpp
```

• ...that we will refactor to our shellcode, and compile by masm:

```
ml <file_name>.asm
```

- it will generate a PE: we will cut out the code section, that is our shellcode
- The key is the refactoring! We need to follow all the principles of building shellcodes...

Creating shellcode: C

• Use the given template, and refactor the application in C into a valid shellcode, by following the steps...

Exercise time...



- From a C project, through assembly, to shellcode:
 - https://vxug.fakedoma.in/papers/VXUG/Exclusive/FromaCprojectthroughassemblytoshellcodeHas herezade.pdf