

Shellcode: advantages

- Self-sufficient: easy to inject into other applications
- Small: can fit into a tiny space i.e. caves between sections
- May be used as a loader: first code injected into an application, that follows to load other modules
- Sometimes (but less often) the full malicious functionality can be implemented as shellcode (i.e. Fobber malware)
- This type of code was popular in the past, virus era: where malware code was added to existing PE files (rather than injected into processes)

Creating shellcode

- In case of PE format we just write a code and don't have to worry how it is loaded: Windows Loader will do it
- It is different when we write shellcode
- We cannot rely on the conviniences provided by PE format and Windows Loader:
 - No sections
 - No Data Directories (imports, relocations)
 - Only code to provide everything we need...

Creating shellcode

| Feature | PE file | shellcode |
|---|--|--|
| Loading | via Windows Loader running new EXE triggers creation of a new process | Custom, simplified must parasite on existing process (i.e. via code injection + thread injection) |
| Composition | Sections with specific access rights, carrying various elements (code, data, resources, etc) | All in one memory area (read,write,execute) |
| Relocation to the load base | Defined by relocation table, applied by Windows Loader | Cutom; position-independent code |
| Access to system API (Imports loading) | Defined by import table, applied by Windows Loader | Custom: retrieving imports via PEB lookup; no IAT, or simplified |

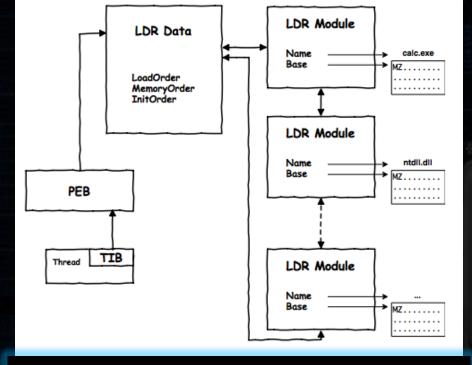
Position-independent code

- In order to create a position-independent code, we must take care that all the addresses that we use are relative to the current intruction pointer address
- A short jump, long jump, call to a local function are relative -> we can use them!

```
EBEO V JMP SHORT 0X413BA8
E8EE000000 V CALL 0X413BFA
```

Any address that needs to be relocated (i.e. using of the data from different PE section)
 breaks the position independence:

• In order to retrieve the imported functions, we will take advantage of the linklist pointed by PEB





http://blog.malcom.pl/2017/shellcode-peb-i-adres-bazowy-modulu-kernel32-dll.html

 In order to retrieve the imported functions, we will take advantage of the linklist pointed by PEB

```
Pyoid Reserved1[12]:
    PPEB ProcessEnvironmentBlock;
    PVOID Reserved2[399];
    BYTE Reserved3[1952];
    PVOID TlsSlots[64];
    BYTE Reserved4[8];
    PVOID Reserved5[26];
    PVOID Reserved5[26];
    PVOID Reserved6[4];
    PVOID TlsExpansionSlots;
} TEB, *PTEB;
```

```
typedef struct PEB
     UCHAR InheritedAddressSpace;
    UCHAR ReadImageFileExecOptions:
     UCHAR BeingDebugged;
    UCHAR BitField:
    ULONG ImageUsesLargePages: 1:
     ULONG IsProtectedProcess: 1:
    ULONG IsLegacyProcess: 1;
    ULONG IsImageDynamicallyRelocated: 1;
    ULONG SpareBits: 4;
     PVOID Mutant:
     PVOID ImageBaseAddress:
    PPEB LDR DATA Ldr;
     FLS CALLBACK INFO * FlsCallback;
    LIST ENTRY FlsListHead;
     PVOID FlsBitmap;
    ULONG FlsBitmapBits[4];
    ULONG FlsHighIndex;
     PVOID WerRegistrationData;
     PVOID WerShipAssertPtr;
 PEB, *PPEB;
```

```
typedef struct _PEB_LDR_DATA

{
    ULONG Length;
    BOOLEAN Initialized;
    HANDLE SSHandle:
    LIST_ENTRY InLoadOrderModuleList;
    LIST_ENTRY InMemoryOrderModuleList;
    LIST_ENTRY InInitializationOrderModuleList;
    PVOID EntryInProgress;
}

PEB_LDR_DATA, *PPEB_LDR_DATA;
```

• We will process each entry, searching for the DLL that we need...

```
typedef struct _PEB_LDR_DATA

{
    ULONG Length;
    BOOLEAN Initialized;
    HANDLE SsHandle;
    LIST_ENTRY InLoadOrderModuleList;
    LIST_ENTRY InMemoryOrderModuleList;
    LIST_ENTRY InInitializationOrderModuleList;
    PVOID EntryInProgress;
}
PEB_LDR_DATA, *PPEB_LDR_DATA;
```

```
lkd> dt ntdll! LDR DATA TABLE ENTRY
   +0x000 InLoadOrderLinks : LIST ENTRY
   +0x010 InMemoryOrderLinks : _LISI_ENIRY
                                     : LIST ENTRY
   +0x020 InInitializationOrderLinks
   +0x030 DllBase
                           : Ptr64 Void
   +0x038 EntryPoint
                             Ptr64 Void
   +0x040 SizeOfImage
                             Uint4B
  +0x048 FullDilName
                              HNICODE STRING
   +0x058 BaseDllName
                              UNICODE STRING
   +0x068 FlagGroup
                               l] UChar
```

Next LDR_DATA_TABLE_ENTRY

typedef struct _UNICODE_STRING {
 USHORT Length;
 USHORT MaximumLength;
 PWSTR Buffer;
} UNICODE_STRING, *PUNICODE_STRING;

L"Ntdll.dll"

- 1. Get the PEB address
- 2. Via PEB->Ldr->InMemoryOrderModuleList, find:
 - kernel32.dll (which is always loaded)
 - or ntdll-dll (if we want to use low-leven equivalents of Import loading functions)
- 3. Walk through exports table to find addresses of:
 - LoadLibraryA (eventually: ntdll.LdrLoadDll)
 - GetProcAddress (eventually: ntdll LdrGetProcedureAddress)
- 4. Use LoadLibrary A to load other needed DLLs
- 5. Use GetProcAddress to retrieve functions

Creating shellcode: assembly

• We can use YASM for shellcodes written in pure assembly:

```
yasm -f bin demo∙asm
```

- We will not use a linker, which means:
 - we need to fill imports by ourselves
 - we need to take care of relocations or make the code position-independent

Creating shellcode: C

• We can use a C compiler to generate assembly:

```
C1 /c /FA <file_name>.cpp
```

• ...that we will refactor to our shellcode, and compile by masm:

```
ml <file_name>.asm
```

- it will generate a PE: we will cut out the code section, that is our shellcode
- The key is the refactoring! We need to follow all the principles of building shellcodes...

Creating shellcode: C

• Use the given template, and refactor the application in C into a valid shellcode, by following the steps...

Exercise time...



- From a C project, through assembly, to shellcode:
 - https://vxug.fakedoma.in/papers/VXUG/Exclusive/FromaCprojectthroughassemblytoshellcodeHas herezade.pdf