## Malware Writeup

## Research

I knew for this project I wanted to do malware analysis. My initial research consisted of finding a malware sample, which I did through github.com/ytisf/theZoo.

## Analysis

This malware is a 32 bit statically linked elf binary meant to run on the linux architecture. I first did some static analysis on the binary, but found this somewhat limiting. Eventually I moved to debug the binary from a bootable disk.

I started my analysis process by looking at the list of functions that radare2 identified. Because the binary is statically linked, there were a lot of extra functions related to system calls, but overall the binary still wasn't too big. The most obvious function to investigate first was the main function.

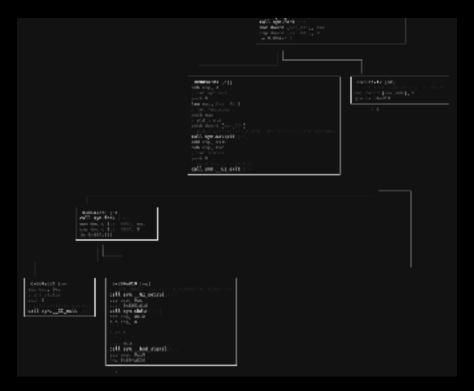
[0x6864a769	]> af	ll-:												
address	size	nbbs	edges	0.0	cost	min bound	range	nex bound	calls	locals	args	xref	frene	nare
0×68648687	58	4	4	- 2	30	0x00040087	58	0x08648641	1		2	- 5	28	sym.printchar
6x68648641	215	26	27	9	92	6x68648641	215	6x68648718	3	4	4	4		sym.prints
6x68648718	293	18	26	19	115	6x68648718	293	6x6864883d	3	9	7	4		sym.printi
6x6864883d	584	33	48	17	248	6x6864883d	584	6x68648a85	7	7	3	1		sym.print
0x0804b16c	27	3	3	2	17	0x0804b16c		6x6864b187	9	ė	2	5		syn.strcpy
0x08048be8	396	13	18	7	133	0x08048be8	396	0x08048d6c	3	14	3	1		sym.recvLine
0x0804b8a4	26	- 1		i	12	0x0804b8a4		0x0804b0be	ī	0	ī	2	28	syn, waitpid
0x6864d43c	59	3	3	Z	29	6x6864d43c	59	0x0864d477	ī	ō	4	ī	16	sym.wait4
oxessaesf4	19			ī	12	0x6964e8f4	19	0x6964e987		o	1	9	- 4	sym.strlen
exene4d554	39				22	exensad554	39	exece4d57b		0	3	1.8		зув. вивору
oxene4e5dc	98	- 5	n	3	52	exens4e5dc	98	0x0864e63e	4	1	2	- 2	44	syn.fgets
exeae4d2da	50	3	3		25	exeas4d2da	50	exeae4d382	1	0	2	1	12	sym.getrlimit
exess4acfc	63	3	3	2	28	exess4acfc	63	exese4af3b	1	1	4	2	32	syn.loctl
6x6864d294	38	3	3	2	21	exese4d294	38	6x6864d2ba	1	0	8	1	12	syn.getgid
exese4ccbc	381	38	36	156	125	exessaccbc	325	6x6864cc81	2	1	2	1	28	syn.sysconf
0x0804d21c	37	3	3	2	18	0x0804d21c	37	0x6864d241	1	2	0	1	44	syn.getdtablestze
0x0804c604	72	- 1	0	ī	34	0x0804c684	72	0x0804c64c		3	0	1	60	sym, random
0x08649719	1355	36	58	15	504	0x68649719	1355	0x08649c64	27	36	7	2	204	syn.atcp
0x6864b735	51			1	19	6x6864b735	51	0x0864b763	1	4	4	2	44	sym.recv
охоночьюче	43		8	1	17	оженечьем:	43	0x6864b6c7	1	3	- 3	- 4	44	sym.connect
0x0864a634	309	11	14		114	0x6864a634	309	0x0864a769	9	3	н	1	556	sym.initConnection
0x08851158	176	19	29	12	86	0x08851158	176	0x08651288			3	1	12	sym.memrchr
exe884d26c	38	3	3		21	exe884d26c	38	exe8e4d292	1	0	8	1	12	sym.geteuid
exe884d57c	39	4	4	2	22	exe884d57c	39	exe8e4d5a3	- 0	0	3	2	8	sym.memmove
6x6864d32c	59	3	3	2	25	6x6864d32c	50	6x6864d35c	1	0	2	1	12	sym.munmap
6x6864c988	29	- 1	8	1	11	6x6864c988	20	6x6864c994	1	0	1	31	28	sym.atol
0x0864c994	26	- 1	0	1	12	0x0804c994	26	0x0804c9ae	1	0	1	1	28	sym.strtol
0x0804d244	38	3	3	2	21	0x0804d244	38	0x0804d26a	1	0		1	12	syn.getegid
0x0864aed4	38	3	3	2	21	0x6864aed4	38	0x0864aefa	1	0			12	syn.getpid
0x68648b3e	59	4	4	- 2	27	6x68648b3e	59	0x68648b79	1		- 2	- 3	44	syn.getHost
охонечиь79	103	4	4		45	ежене <b>4</b> нБ79		Oxene4iibe8	1		2		36	sym.makeRandomStr
0x686482b7	45		н	1	24	ежене4н267	45	0x888482e7			1	4	44	sym.getRandomIP
exe884b2a3			8	1		exe884b2a3		exe8e4b2aa		0	1	1		sym.ntohl
exese4afbb	25		8	1	11	exeae4afbb		exese4afd4	1	0	2	9	28	sym.creat
exe8e4b818	43	1	8	1	17	exe8e4b818	43	6x6864b843	1	3	3	7	44	sym.socket
exe884b258	21	- 1	1	9		6x6864b258		exese4b265	9	Θ	2	1	0	sym.bcopy
6x6864b764	51	- 1	8	1	19	6x6864b764		0x0864b797	1	4	4	3	44	syn. send
0x08648558	175	11	14	5	81	0x08048558		0x08648687	2	5	1	1	44	sym.trim
0x6864d384	38	3	3	2	21	0x0804d384		0x0864d32a	1	0		1	12	syn.getuid
0x6864ea88	29			1	14	0x6864ea58	29	0x0864ea25	1		1	- 3	92	syn.isatty
0x686564a8	2.2		н	1	18	0×696584a8		0x086564c9	1	0	- 2	1	78	зув. виврору
охонвасьза	393	23	32	11	174	ехене4с638	393	exess4cch9	12	9	1	1	460	sym.sleep

Several things happen in the first block of the main function. First, srandom() is called using the current time and the pid of the binary as seeds for it.

Then, init\_rand() is called using the xor of the time and the pid of the binary once again. After this, a function is called that gets the IP address of the local computer.

```
call sym._GI_time;[ca]
add csp, 8x10
nov ebx, esx
. int getped(void)
call sym.getpid;[co]
xor eax, ebx
aub wap, 8xc
; int seed
push eax
; void ar and(int seed)
call sym.srand;[co]
add csp, 8x10
aub wap, 8xc
; time t *timer
push 0
: time t *timer
push 0
```

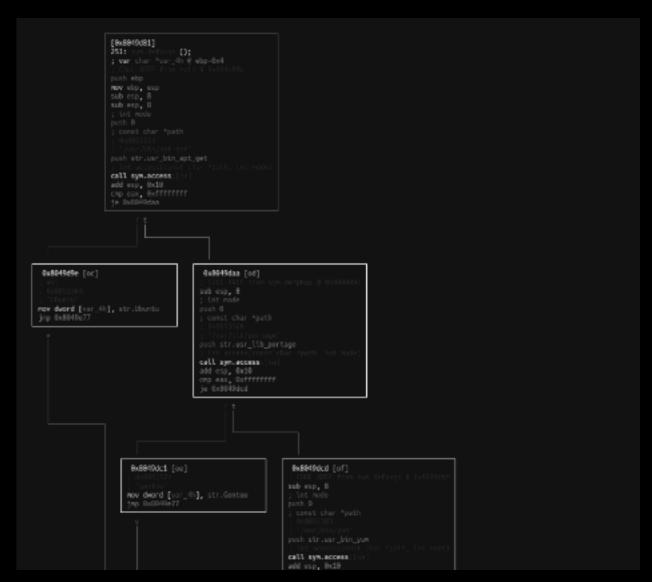
After this the binary forks twice I think as a bad anti-debugging trick. The parent processes will jump to a block that exits the binary and the final child process will continue on with the rest of the execution.



This child process then enters the main execution loop of the binary. The first step is to initialize some things, including changing to the / directory. Then,

the malware checks if there is a server available at 209.141.48.138:666. If this connection fails the binary will sleep for 5 seconds and loop this check again. If it eventually succeeds, it moves on to the next phase of execution.

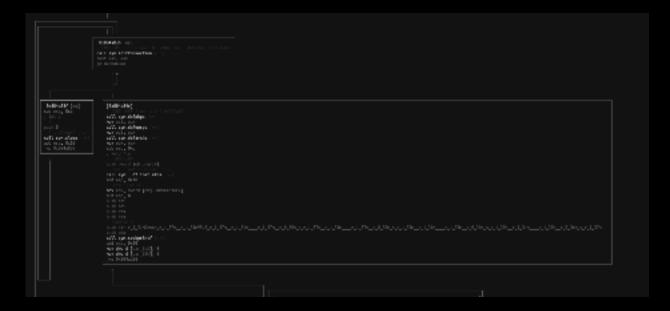
In the next block it gets some basic information about the system it is running on. In the first function call, it checks whether it has access to various directories, including apt-get, /usr/lib/portage, /usr/bin/yum, /usr/share/YaST2, /usr/etc/pkg and some others. Based on these directories it can guess what distribution is running on the machine because the directories correspond to the default package manager of each distribution.



In the next two functions the binary returns that it is running on a linux machine and that it is using the  $x86_32$  architecture. After this there is a call to sockprintf that passes the following string

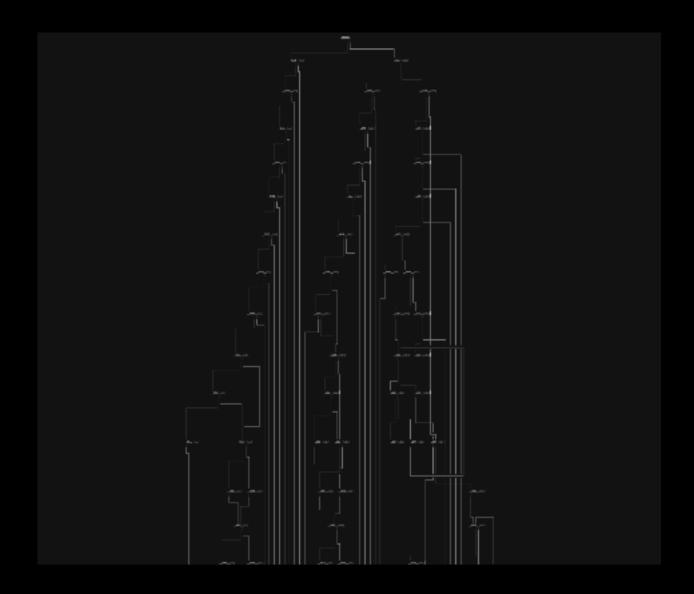
\e[1;31mDemon\e[1;37m[\e[1;31mV5.0\e[1;37m]\e[1;31m-->\e[1;37m[\e[0;36m%s\e[1;37m]\e[1;31m-->\e[1;37m[\e[0;36m%s\e[1;37m]\e[1;31m-->\e[1;37m[\e[0;36m%s\e[1;37m]\e[1;31m-->\e[1;37m]\e[

What this is doing is passing information about the architecture, distro, etc to the central command server.



After this is the main loop that acts as the payload. At the beginning of this loop it attempts to receive a command from the main control server using a readline() function. This input is split up and converted to uppercase characters. These strings are then passed to the cncinput() function.

This function contains the core payload functionality. It uses the command information coming from the central server to decide which behavior to branch to. The most notable of these behaviors is a DOS attack. This malware does not self propagate. As you can see below, the control flow is complicated.



Analysis revealed the following command that could be used by the malicious server to launch various DDOS attacks.

## **Command syntax**

- 1. UDP <ip> <port> <time> <spoofit> <packetsize> <pollinterval>
  - a. UDP DDoS attack with random payload
  - b. Time of attack in seconds
  - c. If no port specified, random port is generated every <pollinterval>
     packets
  - d. Random data does not change during attack
- 2. TCP <ip> <port> <time> <spoofit> <flags> <packetsize> <pollint>
  - a. TCP DDoS attack

- b. If IP is id, TCP seq and TCP source port change randomly every <pollinterval> packets
- c. Payload does not change during attack but changes bot to bot
- d. Flags include all or syn,rst,fin,ack,psh
- 3. STD <ip> <port> <time>
  - a. UDP attack with fixed payload
- 4. STOMP <ip> <port> <time> <spoofit> <flags> <packetsize> <poolint>
  - a. STD attack followed by UDP attack followed by TCP attack
- 5. CNC <ip> <port> <time>
  - a. Make TCP connection to server with <ip> on port <port> and closes after
     1 second, repeating continually
- 6. STOP
  - a. Stop all attacks. Each attack gets it's own fork/process. Main process keeps track of all the forks and kills them all if this command is called.