I nsea hon_sout

we take an element, compare its value with left-side element,

if subsequent left side element is greater than selected element,

we shift the element to right, we will compare and shift right until any left subsequent element is less than selected element on we have reach zeroth element, and then insert the element at that position.

we start to loop through (pass through) from 1st index.

we will store the index position and element value.

position = index

temp_value = arr [index]

we will compare with left subsequent value.

4 2 7 1 3

position = 1 $temp_value = 2$

is 4 > 2 ? Yes

then, shift '4' to right, and update the position.

4 4 7 1 3

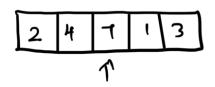
position = 0 $temp_val = 2$

Here we have reached left most position and nothing to compare on left side, So insent the temp-val at the position.

2 4 7 1 3

we can say that these element are sorted in order.

NOW In 2nd passith sough,



position = 2, $temp_val = 7$

we will compose the left subsequent value

is H > 7? No, Hence pass-through ends



1 —) after 2rd pass+hoovegh, we can say elements one souted till 2rd index.

In 3rd pasthrough,

position = 3 temp_val = 1

is 7 > 1 ? Yes

we will shift 7 to right, and update the position.

position = 2 temp-val = 1

10 4>1 7 Yes

snift 4 to sight, update position.

basinson = 1

temp_val = 1

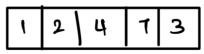
is 2 > 1 ? Y2

Suift 2 to right, update position.

position = 0

temp_val = 1

Nothing to compare on left, Hence ineest temp_rol in the position.



—) after 3rd pasthrough, we can say, elements are sorted till 3rd index.

In 4th pasthrough,

position = 4

temp_val = 3

18 7 > 37 Yes Shift I to right, update position

| 1 | 2 | 4 | 7 | 7 | position = 3

is 4>3? Yes

shift 4 to right, update position

1 2 4 4 7

position = 2

temp-val = 3

is 2 > 3? No, so we will insert temp-val at position.

12347

so, after 4 pass through we can say away is sould.

- ·- For N elements, it will take (N-1)
 nack through
- .. After each passthrough, we can say elements till passthrough index are sorted.

Implementation:

fune Inspirion Soprit (com)

for (index = 1 \rightarrow len(arr)-1) {

notition = index

temp_val = const [index]

while (position > 0 27

are [position] > temp-val)

are [position] = are [position-1]

position--:
3

ar [position] = temp_val;

z

Efficiency:

There are 3 types of steps that occur, compasison, shifting 7 Insertion.

- (1). comparison In worst case scenario, we will be comparing all the left-side elements $\therefore \frac{N(N-1)}{2} \text{ Steps.}$
 - (11). Shifting in worst case, we will be shifting all the left-side element

to right for each passtmough

(111) Insertion - for each shift, insertion nappens tuice

$$= N^2 - N + N^2 - N + 2N - 2$$

=>
$$N^2 - N + N^2 - N + 2N - 2$$

=> $N^2 - 2$ (appointinately)

: Insertion sort is Big O(N2).

Another major rule.

Big O Notation takes into account the highest order of N.

eq:
$$N^4 + N^2 + N^2 + N$$

The average case:

- In worst case: selection sort > Insertion sort
 - The cases that occurs most frequently are overage scenario.
 - worst and best case occurs manely.
 - Insertion suset
 - · worst cose : O(N2)
 - · Average-case : O(N2)
 - · Best-case: O(n)
 - Selection sout

 - · wort coxe · Average cose · Best case

- Selection Sout _ Don't have any mechanism to end passthoogh early at any point, Fach mother ugh

- compares value to its right no matter what.
- Insertion sout when the values
 of left -side is souted, no shifting
 is done, thus shifting steps are
 seduced.
- → what is best?
 It depends,
 - when array is randomly sorted then insertion sort.
 - when away is souted in reverse order than selection sout.
 - when no idea, how data will be present, then on average case both will work fine.