Selection_sout

Civen an array, we find the lowest element value and swap it in order.

Steps

- we loop through away, from left to right.
- we keep a track of the lowest_index_val.
- for each pass through, if lowest value is not in position, then we swap the passthrough index with lowest value.
- We repeat the above steps till the second last index is souted.

eg: 0 1 2 3 4 4 2 7 1 3

Ist pass through,

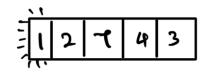
Initially, assuming the first element to be least, so lowest_index_value has the '0' index.

we loop through arriay, and find the lowest value and if found we will update the lowest-index-value.

Here in this case, index '3' is the lowest, we will compare it with the

posstrinough index ('0' in this case, outer loop index), if both one different, then swap the value.

so now the areay becomes,



we know that, first index of the armay is Sorted now, so 2^{nd} pass through begins from 2^{nd} index (index 1).

we loop through and find the lowest_index_value, if both are different we swap.

So after and past through, curray becomes

Since no value was lesser than the lowest-index-value and lowest-index-value and lowest-index-value was equal to passthrough index, swapping isn't done, exemend is in convect order.

In 3rd pass+snrough, pass+snrough index is '2', assuming it to be the lowest value, booking through the remaining blements of

tro overy, finding the lowest and swapping.
So assnay becomes,

in when pass through, we follow the above steps, after which, and last element is sorted,

Since all the Rlements till 2nd last has been would, we can say that last element is also sould, and end the program.

Implementation:

func selectionSort (ax)

for (i = 0 -> len(arr) -1)

lowest_ind_val = i

Finding lowest vellue Indux

for
$$(j=i+1 \rightarrow len(\omega x)-1)$$

if $(\alpha r[j] \perp \alpha r[lowest_ind_val])$
 $lowest_ind_val = j$

if (lowest_ind_val 1= i)

swap (ar, lowelt_ind_val, i)

refusin cror;

Efficiency:

Selection sout how two kinds of steps,

- (1). Comparison where we compare and find the lowest index value.
- (11). Swap Swap the element, swapping is done either ones or zero, for each passthrough.

for N ecements,

(1). Companison is done for $\frac{N(N-1)}{2}$ steps.

(11). Swapping is done, either ones on 2000.

Since Swapping is done maximum one por possition of lowest-index-value is not in consect position else no swapping is done. Compare to Bubble sout, in worst-case, we have to swap for each comparison.

This makes selection sout faster than

Bubble sort

mathamounically, Selection Sout takes $\left(\frac{N_2^2}{2}\right)$ steps in woust case and Bubble sout takes $\left(\frac{N^2}{2}\right)$.

Selection sort is twice as fast as Bubble sort but one of the major rule of Big 0 is that Ignore Constants, because of which both selection sort and bubble sort because $O(n^2)$.

Big O Notation Ignores Constants

Pole of Big 0:

purpose of Big O is that for different clossification, there will be a point at which one classification supersedes the other in speed, and will remain faster forever. When that point occurs exactly, However, it is not the concern of Big O.

It two algorithm falls under 2 different classification, we can easily say which allower with large

amount of date.

However, when it falls under some classification, further analysis is required to determine which algorithm is faster.

As falling under the same classification doesn't necessarily means both process the same speed.

eg: selection sort and Bubble sout.