

# You could've invented Strassen's Algorithm

Zeke

July 2025

# Matrix Multiplication

Given two  $2 \times 2$  matrices  $A, B$

$$A = \begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix}, \quad B = \begin{pmatrix} b_{11} & b_{12} \\ b_{21} & b_{22} \end{pmatrix}$$

# Matrix Multiplication

Given two  $2 \times 2$  matrices  $A, B$

$$A = \begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix}, \quad B = \begin{pmatrix} b_{11} & b_{12} \\ b_{21} & b_{22} \end{pmatrix}$$

Their product  $C = AB$  is

$$C = \begin{pmatrix} a_{11}b_{11} + a_{12}b_{21} & a_{11}b_{12} + a_{12}b_{22} \\ a_{21}b_{11} + a_{22}b_{21} & a_{21}b_{12} + a_{22}b_{22} \end{pmatrix}$$

# Matrix Multiplication

Given two  $2 \times 2$  matrices  $A, B$

$$A = \begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix}, \quad B = \begin{pmatrix} b_{11} & b_{12} \\ b_{21} & b_{22} \end{pmatrix}$$

Their product  $C = AB$  is

$$C = \begin{pmatrix} a_{11}b_{11} + a_{12}b_{21} & a_{11}b_{12} + a_{12}b_{22} \\ a_{21}b_{11} + a_{22}b_{21} & a_{21}b_{12} + a_{22}b_{22} \end{pmatrix}$$

But computing  $C$  requires 8 multiplications. Can we do better?

# Strassen's Algorithm

Remarkably, yes.  $C$  can be computed with 7 multiplications.

$$P_1 = (a_{11} + a_{22})(b_{11} + b_{22})$$

$$P_2 = (a_{21} + a_{22})b_{11}$$

$$P_3 = a_{11}(b_{12} - b_{22})$$

$$P_4 = a_{22}(b_{21} - b_{11})$$

$$P_5 = (a_{11} + a_{12})b_{22}$$

$$P_6 = (a_{21} - a_{11})(b_{11} + b_{12})$$

$$P_7 = (a_{12} - a_{22})(b_{21} + b_{22})$$

$$C = \begin{pmatrix} P_1 + P_4 - P_5 + P_7 & P_3 + P_5 \\ P_2 + P_4 & P_1 - P_2 + P_3 + P_6 \end{pmatrix}$$

# My Hackathon Project

- But this is arcane...

$$C = \begin{pmatrix} P_1 + P_4 - P_5 + P_7 & P_3 + P_5 \\ P_2 + P_4 & P_1 - P_2 + P_3 + P_6 \end{pmatrix}$$

# My Hackathon Project

- But this is arcane...

$$C = \begin{pmatrix} P_1 + P_4 - P_5 + P_7 & P_3 + P_5 \\ P_2 + P_4 & P_1 - P_2 + P_3 + P_6 \end{pmatrix}$$

- What if we could automatically discover this with a computer program?

# My Hackathon Project

- But this is arcane...

$$C = \begin{pmatrix} P_1 + P_4 - P_5 + P_7 & P_3 + P_5 \\ P_2 + P_4 & P_1 - P_2 + P_3 + P_6 \end{pmatrix}$$

- What if we could automatically discover this with a computer program?
- Would we learn something about optimizing machine learning workloads along the way?



# An algebra problem

1. Let  $\bar{a}$  be  $(a_{11}, a_{12}, a_{21}, a_{22})$  and  $\bar{b}, \bar{c}$  be likewise.
2. Neither  $a_{ij}a_{kl}$  nor  $b_{ij}b_{kl}$  appears in  $C$ . So, the  $k$ th product will be of the form

$$P_k = \left( \sum_i \bar{a}_i \mu_i^{(k)} \right) \left( \sum_j \bar{b}_j \varphi_j^{(k)} \right)$$

where  $\mu_i^{(k)}, \varphi_j^{(k)} \in \{0, 1\}$  are 1 if and only if  $\bar{a}_i/\bar{b}_j$  appear in the  $k$ th product.

# An algebra problem

1.  $\bar{c}_l$  will be a combination of  $P_k$ s.

# An algebra problem

1.  $\bar{c}_I$  will be a combination of  $P_k$ s.
2. Letting  $w_I^{(k)} \in \{-1, 0, 1\}$  denote  $P_k$ 's coefficient in the expression for  $\bar{c}_I$ ,

$$\bar{c}_I = \sum_k w_I^{(k)} P_k$$

## Putting it together

$$\begin{aligned}\bar{c}_l &= \sum_k w_l^{(k)} P_k \\ &= \sum_k w_l^{(k)} \left( \sum_i \bar{a}_i \mu_i^{(k)} \right) \left( \sum_j \bar{b}_j \varphi_j^{(k)} \right) \\ &= \sum_k w_l^{(k)} \sum_i \sum_j \bar{a}_i \bar{b}_j \mu_i^{(k)} \varphi_j^{(k)} \\ &= \sum_i \sum_j \bar{a}_i \bar{b}_j \sum_k w_l^{(k)} \mu_i^{(k)} \varphi_j^{(k)}\end{aligned}$$

## 8 multiply equivalence

1. To make this an algebra problem, we need a formula for  $\bar{c}_l$ .

## 8 multiply equivalence

1. To make this an algebra problem, we need a formula for  $\bar{c}_l$ .
2. Recall,  $\bar{c}$  is a flattened version of

$$C = \begin{pmatrix} \bar{a}_1 \bar{b}_1 + \bar{a}_2 \bar{b}_3 & \bar{a}_1 \bar{b}_2 + \bar{a}_2 \bar{b}_4 \\ \bar{a}_3 \bar{b}_1 + \bar{a}_4 \bar{b}_3 & \bar{a}_3 \bar{b}_2 + \bar{a}_4 \bar{b}_4 \end{pmatrix}$$

## 8 multiply equivalence

1. To make this an algebra problem, we need a formula for  $\bar{c}_l$ .
2. Recall,  $\bar{c}$  is a flattened version of

$$C = \begin{pmatrix} \bar{a}_1 \bar{b}_1 + \bar{a}_2 \bar{b}_3 & \bar{a}_1 \bar{b}_2 + \bar{a}_2 \bar{b}_4 \\ \bar{a}_3 \bar{b}_1 + \bar{a}_4 \bar{b}_3 & \bar{a}_3 \bar{b}_2 + \bar{a}_4 \bar{b}_4 \end{pmatrix}$$

3. Letting  $T_{ijl}$  be 1 if  $\bar{a}_i \bar{b}_j$  appears in the above expression for  $\bar{c}_l$  and 0 otherwise gives a formula for  $\bar{c}_l$  as desired.

$$\bar{c}_l = \sum_i \sum_j \bar{a}_i \bar{b}_j T_{ijl}$$

# Putting it together

Because

$$\bar{c}_l = \sum_i \sum_j \bar{a}_i \bar{b}_j \sum_k w_l^{(k)} \mu_i^{(k)} \varphi_j^{(k)}$$

and

$$\bar{c}_l = \sum_i \sum_j \bar{a}_i \bar{b}_j T_{ijl}$$

we have

$$T_{ijl} = \sum_k w_l^{(k)} \mu_i^{(k)} \varphi_j^{(k)}$$



# Putting it together

Because

$$\bar{c}_l = \sum_i \sum_j \bar{a}_i \bar{b}_j \sum_k w_l^{(k)} \mu_i^{(k)} \varphi_j^{(k)}$$

and

$$\bar{c}_l = \sum_i \sum_j \bar{a}_i \bar{b}_j T_{ijl}$$

we have

$$T_{ijl} = \sum_k w_l^{(k)} \mu_i^{(k)} \varphi_j^{(k)}$$

an algebra problem, the solution of which corresponds to a way to compute  $C = AB$  using  $k$  multiplies.

# Conclusion

1. If you put this into a computer program for solving algebra problems, it finds Strassen's algorithm. :)
2. Have we learned something deep about optimizing machine learning workloads?

# Conclusion

1. If you put this into a computer program for solving algebra problems, it finds Strassen's algorithm. :)
2. Have we learned something deep about optimizing machine learning workloads?  $\neg\_(\_)\_/$

# Conclusion

1. If you put this into a computer program for solving algebra problems, it finds Strassen's algorithm. :)
2. Have we learned something deep about optimizing machine learning workloads?  $\sim \backslash (\text{シ}) \_ / \sim$
3. Thanks for having me!