Recursive Descent Parser

1. Recursive Descent Parser: It is a kind of Top-Down Parser. A top-down parser builds the parse tree from the top to down, starting with the start non-terminal. By carefully writing a grammar means eliminating left recursion and left factoring from it, the resulting grammar will be a grammar that can be parsed by a recursive descent parser.

Rule for immediate left recursion removal.

- For a rule like:
 - A -> A α | β
- · We may replace it with
 - A → β A'
 - A' -> α A' |ε
- (a) Write a program to remove the recursion from the following grammars:

```
Grammar 1 : E \rightarrow E+T|T; T \rightarrow T*F|F; F \rightarrow (E)|id 2 marks Grammar 2 : S \rightarrow Aa | b; A \rightarrow Ac | Sd | \epsilon 3 marks
```

- (b) Recursive Descent Parser:
 - · Consists of a set of procedures, one for each nonterminal
 - Execution begins with the procedure for start symbol
 - A typical procedure for a non-terminal A is gien by:

```
void A() {
   choose an A-production, A->X1X2..Xk
   for (i=1 to k) {
      if (Xi is a nonterminal)
          call procedure Xi();
      else if (Xi equals the current input symbol a)
          advance the input to the next symbol;
      else /* an error has occurred */
   }
}
```

Implement a recursive descent parser and show the results (Whether the string is valid or not and the parse tree generated) for the grammar 1 above and test for input strings: (id+id)*id,)id+id*id 5 marks