### 清华大学 代码册

# 1、几 25

1.1	注意	25
1.2	几何公式	25
1.3	多边形	27
1.4	多边形切割	
1.5	浮点函数	31
1.6	面积	36
1.7	球面	37
1.8	三角形	38
1.9	三维几何	40
1.10	凸包	47
1.11	网格	49
1.12	圆	49
1.13	整数函数	51

# 2、组 54

2.1	组合公式	54
2.2	排列组合生成	54
2.3	生成 gray 码	56
	置换(polya)	
	字典序全排列	
2.6	字典序组合	57

# 3、结 58

3.1	并查集	58
3.2	堆	59
3.3	线段树	60
3.4	子段和	65
3.5	子阵和	65

# **4、数 66**

4.1	阶乘最后非0位	66
4.2	模线性方程组	67
4.3	素数	68
4.4	欧拉函数	69

# 5、数值计算

5.1	定积分计算(Romberg)	70
	多项式求根(牛顿法)	
5 3	周期性方程(追赶法)	73

### 6、图论一NP搜索

6.1	最大团	74
6.2	最大团(n<64)(faster)	75

### 7、图论一连通性

7.1	无向图关键点(dfs 邻接阵)	77
7.2	无向图关键边(dfs 邻接阵)	78
7.3	无向图的块(bfs 邻接阵)	79
7.4	无向图连通分支(dfs/bfs 邻接阵)	80
7.5	有向图强连通分支(dfs/bfs 邻接阵)	81
7.6	有向图最小点基(邻接阵)	82

## 8、图论一匹配

8.1 二分图最大匹配(hungary 邻接表)	83
8.2 二分图最大匹配(hungary 邻接阵)	84
8.3 二分图最大匹配(hungary 正向表)	84
8.4 二分图最佳匹配(kuhn_munkras 邻接阵)	85
8.5 一般图匹配(邻接表)	86
8.6 一般图匹配(邻接阵)	87
8.7 一般图匹配(正向表)	

## 9、图论一网络流

9.1	最大流(邻接阵)	88
	上下界最大流(邻接阵)	
	上下界最小流(邻接阵)	
9.4	最大流无流量(邻接阵)	91
9.5	最小费用最大流(邻接阵)	91

# 10、图论一应用 92

10.1	欧拉回路(邻接
阵)	92

## 10.2 树的前序表转化

### 10.3 树的优化算法94

10.4	拓扑排序(邻接
阵)	95

### 10.5 最佳边割集...96

### 10.6 最佳点割集...97

### 10.7 最小边割集...98

### 10.8 最小点割集...99

# 10.9 最小路径覆盖 101

### 11、图论一支撑树

11.1	最小生成树(kruskal 邻接表)	101
11.2	最小生成树(kruskal 正向表)	103
11.3	最小生成树(prim+binary_heap 邻接表)	104
11.4	最小生成树(prim+binary_heap 正向表)	105
11.5	最小生成树(prim+mapped_heap 邻接表)	106
11.6	最小生成树(prim+mapped_heap 正向表)	108
11.7	最小生成树(prim 邻接阵)	109
11.8	最小树形图(邻接阵)	109

### 12、图论一最短路径

12.1	最短路径(单源	bellman_ford 邻接阵)	111
12.2	最短路径(单源	dijkstra+bfs 邻接表)	111
12.3	最短路径(单源	dijkstra+bfs 正向表)	112
12.4	最短路径(单源	dijkstra+binary_heap 邻接表)	113
12.5	最短路径(单源	dijkstra+binary_heap 正向表)	114
12.6	最短路径(单源	dijkstra+mapped_heap 邻接表)	115
12.7	最短路径(单源	dijkstra+mapped_heap 正向表)	116
12.8	最短路径(单源	dijkstra 邻接阵)	117
		floyd warshall 邻接阵)	

# 13、应

### 用

13.1 Joseph 问题	118
13.2 N 皇后构造解	119
13.3 布尔母函数	120
13.4 第 k 元素	120
13.5 幻方构造	121
13.6 模式匹配(kmp)	122
13.7 逆序对数	123
13.8 字符串最小表示	123
13.9 最长公共单调子序列	124
13.10 最长子序列	125
13.11 最大子串匹配	126
13.12 最大子段和	127
13.13 最大子阵和	127

## 14、其

### 它

14.1	大数(只能处理正数)	128
14.2	分数	134
14.3	矩阵	136
14.4	线性方程组	138
14.5	线性相关	140
	日期	

### 1、几何

### 1.1 注意

- 1. 注意舍入方式(0.5 的舍入方向);防止输出-0.
- 2. 几何题注意多测试不对称数据.
- 3. 整数几何注意 xmult 和 dmult 是否会出界; 符点几何注意 eps 的使用.
- 4. 避免使用斜率:注意除数是否会为0.
- 5. 公式一定要化简后再代入.
- 6. 判断同一个 2\*PI 域内两角度差应该是 abs(a1-a2)<br/>beta||abs(a1-a2)>pi+pi-beta; 相等应该是 abs(a1-a2)<eps||abs(a1-a2)>pi+pi-eps;
- 7. 需要的话尽量使用 atan2,注意:atan2(0,0)=0, atan2(1,0)=pi/2,atan2(-1,0)=-pi/2,atan2(0,1)=0,atan2(0,-1)=pi.
- 8. cross product =  $|\mathbf{u}|^*|\mathbf{v}|^*\sin(a)$ dot product =  $|\mathbf{u}|^*|\mathbf{v}|^*\cos(a)$
- 9. (P1-P0)x(P2-P0)结果的意义:

正: <P0,P1>在<P0,P2>顺时针(0,pi)内 负: <P0,P1>在<P0,P2>逆时针(0,pi)内 0: <P0,P1>,<P0,P2>共线,夹角为 0 或 pi

10. 误差限缺省使用 1e-8!

### 1.2 几何公式

### 三角形:

- 1. 半周长 P=(a+b+c)/2
- 2. 面积 S=aHa/2=absin(C)/2=sqrt(P(P-a)(P-b)(P-c))
- 3. 中线 Ma=sqrt(2(b^2+c^2)-a^2)/2=sqrt(b^2+c^2+2bccos(A))/2
- 4. 角平分线 Ta=sqrt(bc((b+c)^2-a^2))/(b+c)=2bccos(A/2)/(b+c)
- 5. 高线 Ha=bsin(C)=csin(B)=sqrt(b^2-((a^2+b^2-c^2)/(2a))^2)

- 6. 内切圆半径 r=S/P=asin(B/2)sin(C/2)/sin((B+C)/2) =4Rsin(A/2)sin(B/2)sin(C/2)=sqrt((P-a)(P-b)(P-c)/P) =Ptan(A/2)tan(B/2)tan(C/2)
- 7. 外接圆半径 R=abc/(4S)=a/(2sin(A))=b/(2sin(B))=c/(2sin(C))

### 四边形:

D1,D2 为对角线,M 对角线中点连线,A 为对角线夹角

- 1.  $a^2+b^2+c^2+d^2=D1^2+D2^2+4M^2$
- 2. S=D1D2sin(A)/2

(以下对圆的内接四边形)

- 3. ac+bd=D1D2
- 4. S=sqrt((P-a)(P-b)(P-c)(P-d)),P 为半周长

### 正 n 边形:

- R 为外接圆半径,r 为内切圆半径
- 1. 中心角 A=2PI/n
- 2. 内角 C=(n-2)PI/n
- 3. 边长 a=2sqrt(R^2-r^2)=2Rsin(A/2)=2rtan(A/2)
- 4. 面积 S=nar/2=nr^2tan(A/2)=nR^2sin(A)/2=na^2/(4tan(A/2))

### 圆:

- 1. 弧长 l=rA
- 2. 弦长 a=2sqrt(2hr-h^2)=2rsin(A/2)
- 3. 弓形高 h=r-sqrt(r^2-a^2/4)=r(1-cos(A/2))=atan(A/4)/2
- 4. 扇形面积 S1=rl/2=r^2A/2
- 5. 弓形面积 S2=(rl-a(r-h))/2=r^2(A-sin(A))/2

### 棱柱:

- 1. 体积 V=Ah,A 为底面积,h 为高
- 2. 侧面积 S=lp,l 为棱长,p 为直截面周长
- 3. 全面积 T=S+2A

### 棱锥:

- 1. 体积 V=Ah/3,A 为底面积,h 为高 (以下对正棱锥)
- 2. 侧面积 S=lp/2,l 为斜高,p 为底面周长
- 3. 全面积 T=S+A

### 棱台:

- 1. 体积 V=(A1+A2+sqrt(A1A2))h/3,A1.A2 为上下底面积,h 为高(以下为正棱台)
- 2. 侧面积 S=(p1+p2)l/2,p1.p2 为上下底面周长,1 为斜高
- 3. 全面积 T=S+A1+A2

### 圆柱:

- 1. 侧面积 S=2PIrh
- 2. 全面积 T=2PIr(h+r)
- 3. 体积 V=PIr^2h

### 圆锥:

- 1. 母线 l=sqrt(h^2+r^2)
- 2. 侧面积 S=PIrl
- 3. 全面积 T=PIr(l+r)
- 4. 体积 V=PIr^2h/3

### 圆台:

- 1. 母线 l=sqrt(h^2+(r1-r2)^2)
- 2. 侧面积 S=PI(r1+r2)l
- 3. 全面积 T=PIr1(l+r1)+PIr2(l+r2)
- 4. 体积 V=PI(r1^2+r2^2+r1r2)h/3

### 球:

- 1. 全面积 T=4PIr^2
- 2. 体积 V=4PIr^3/3

### 球台:

- 1. 侧面积 S=2PIrh
- 2. 全面积 T=PI(2rh+r1^2+r2^2)
- 3. 体积 V=PIh(3(r1^2+r2^2)+h^2)/6

### 球扇形:

- 1. 全面积 T=PIr(2h+r0),h 为球冠高,r0 为球冠底面半径
- 2. 体积 V=2PIr^2h/3

### 1.3 多边形

```
#include <stdlib.h>
#include <math.h>
#define MAXN 1000
#define offset 10000
#define eps 1e-8
#define zero(x) (((x)>0?(x):-(x))<eps)
#define _sign(x) ((x)>eps?1:((x)<-eps?2:0))
struct point {double x,y;};
struct line {point a,b;};

double xmult(point p1,point p2,point p0) {
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);</pre>
```

```
}
//判定凸多边形,顶点按顺时针或逆时针给出,允许相邻边共线
int is convex(int n,point* p){
    int i,s[3]=\{1,1,1\};
    for (i=0;i \le n\&\&s[1]|s[2];i++)
        s[ sign(xmult(p[(i+1)\%n],p[(i+2)\%n],p[i]))]=0;
    return s[1]|s[2];
}
//判定凸多边形,顶点按顺时针或逆时针给出,不允许相邻边共线
int is convex v2(int n,point* p){
    int i,s[3]=\{1,1,1\};
    for (i=0;i \le n\&\&s[0]\&\&s[1]|s[2];i++)
        s[sign(xmult(p[(i+1)%n],p[(i+2)%n],p[i]))]=0;
    return s[0]&&s[1]|s[2];
}
//判点在凸多边形内或多边形边上,顶点按顺时针或逆时针给出
int inside convex(point q,int n,point* p){
    int i,s[3]=\{1,1,1\};
    for (i=0;i \le n\&\&s[1]|s[2];i++)
        s[ sign(xmult(p[(i+1)\%n],q,p[i]))]=0;
    return s[1]|s[2];
}
//判点在凸多边形内,顶点按顺时针或逆时针给出,在多边形边上返回 0
int inside convex v2(point q,int n,point* p){
    int i,s[3]=\{1,1,1\};
    for (i=0;i \le n\&\&s[0]\&\&s[1]|s[2];i++)
        s[sign(xmult(p[(i+1)\%n],q,p[i]))]=0;
    return s[0]\&\&s[1]|s[2];
}
//判点在任意多边形内,顶点按顺时针或逆时针给出
//on edge 表示点在多边形边上时的返回值,offset 为多边形坐标上限
int inside polygon(point q,int n,point* p,int on edge=1){
    point q2;
    int i=0,count;
    while (i<n)
        for (count=i=0,q2.x=rand()+offset,q2.y=rand()+offset;i<n;i++)
             if
(zero(xmult(q,p[i],p[(i+1)\%n]))\&\&(p[i].x-q.x)*(p[(i+1)\%n].x-q.x) < eps\&\&(p[i].y-q.y)*(p[(i+1)\%n])
n].y-q.y) < eps)
```

```
return on edge;
                                          else if (zero(xmult(q,q2,p[i])))
                                                       break;
                                          else
                                                                                                                                                                                                                                                                                 if
(xmult(q,p[i],q2)*xmult(q,p[(i+1)%n],q2) < -eps\&&xmult(p[i],q,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1
%n])<-eps)
                                                        count++;
              return count&1;
}
inline int opposite side(point p1, point p2, point 11, point 12){
              return xmult(11,p1,12)*xmult(11,p2,12)<-eps;
}
inline int dot online in(point p,point 11,point 12){
              return zero(xmult(p,11,12))&&(11.x-p.x)*(12.x-p.x)<eps&&(11.y-p.y)*(12.y-p.y)<eps;
}
//判线段在任意多边形内,顶点按顺时针或逆时针给出,与边界相交返回1
int inside polygon(point 11, point 12, int n, point* p){
              point t[MAXN],tt;
              int i,j,k=0;
              if (!inside polygon(11,n,p)||!inside polygon(12,n,p))
                            return 0;
              for (i=0;i< n;i++)
                            if (opposite side(11,12,p[i],p[(i+1)\%n])&&opposite side(p[i],p[(i+1)\%n],11,12))
                            else if (dot online in(11,p[i],p[(i+1)\%n]))
                                          t[k++]=11;
                            else if (dot online in(12,p[i],p[(i+1)\%n]))
                                          t[k++]=12;
                            else if (dot online in(p[i],11,12))
                                          t[k++]=p[i];
              for (i=0;i<k;i++)
                            for (j=i+1;j< k;j++){
                                          tt.x = (t[i].x + t[j].x)/2;
                                          tt.y=(t[i].y+t[i].y)/2;
                                         if (!inside_polygon(tt,n,p))
                                                       return 0;
                            }
              return 1;
}
point intersection(line u,line v){
```

```
point ret=u.a;
     double t=((u.a.x-v.a.x)*(v.a.y-v.b.y)-(u.a.y-v.a.y)*(v.a.x-v.b.x))
               /((u.a.x-u.b.x)*(v.a.y-v.b.y)-(u.a.y-u.b.y)*(v.a.x-v.b.x));
     ret.x+=(u.b.x-u.a.x)*t;
     ret.y+=(u.b.y-u.a.y)*t;
     return ret;
}
point barycenter(point a,point b,point c){
     line u,v;
     u.a.x=(a.x+b.x)/2;
     u.a.y=(a.y+b.y)/2;
     u.b=c;
     v.a.x = (a.x + c.x)/2;
     v.a.y=(a.y+c.y)/2;
     v.b=b;
     return intersection(u,v);
}
//多边形重心
point barycenter(int n,point* p){
     point ret,t;
     double t1=0,t2;
     int i;
     ret.x=ret.y=0;
     for (i=1;i< n-1;i++)
          if (fabs(t2=xmult(p[0],p[i],p[i+1]))>eps)\{\\
               t=barycenter(p[0],p[i],p[i+1]);
               ret.x+=t.x*t2;
               ret.y+=t.y*t2;
               t1+=t2;
     if (fabs(t1)>eps)
          ret.x/=t1,ret.y/=t1;
     return ret;
}
```

### 1.4 多边形切割

```
//多边形切割
//可用于半平面交
#define MAXN 100
#define eps 1e-8
#define zero(x) (((x)>0?(x):-(x))<eps)
```

```
struct point{double x,y;};
double xmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
int same side(point p1,point p2,point 11,point 12){
    return xmult(11,p1,12)*xmult(11,p2,12)>eps;
}
point intersection(point u1,point u2,point v1,point v2){
    point ret=u1;
    double t=((u1.x-v1.x)*(v1.y-v2.y)-(u1.y-v1.y)*(v1.x-v2.x))
              /((u1.x-u2.x)*(v1.y-v2.y)-(u1.y-u2.y)*(v1.x-v2.x));
    ret.x+=(u2.x-u1.x)*t;
    ret.y+=(u2.y-u1.y)*t;
    return ret;
}
//将多边形沿 11,12 确定的直线切割在 side 侧切割,保证 11,12,side 不共线
void polygon cut(int& n,point* p,point 11,point 12,point side){
    point pp[100];
    int m=0,i;
    for (i=0;i< n;i++)
         if (same side(p[i],side,l1,l2))
              pp[m++]=p[i];
         if
(!same side(p[i],p[(i+1)\%n],11,12)\&\&!(zero(xmult(p[i],11,12))\&\&zero(xmult(p[(i+1)\%n],11,12))))
              pp[m++]=intersection(p[i],p[(i+1)\%n],l1,l2);
    for (n=i=0;i<m;i++)
         if(!i||!zero(pp[i].x-pp[i-1].x)||!zero(pp[i].y-pp[i-1].y))
              p[n++]=pp[i];
    if (zero(p[n-1].x-p[0].x)&&zero(p[n-1].y-p[0].y))
         n--;
    if (n<3)
         n=0;
}
```

### 1.5 浮点函数

```
//浮点几何函数库
#include <math.h>
#define eps 1e-8
```

```
#define zero(x) (((x)>0?(x):-(x))<eps)
struct point{double x,y;};
struct line {point a,b;};
//计算 cross product (P1-P0)x(P2-P0)
double xmult(point p1,point p2,point p0){
     return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
double xmult(double x1,double y1,double x2,double y2,double x0,double y0){
    return (x1-x0)*(y2-y0)-(x2-x0)*(y1-y0);
}
//计算 dot product (P1-P0).(P2-P0)
double dmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.x-p0.x)+(p1.y-p0.y)*(p2.y-p0.y);
double dmult(double x1,double y1,double x2,double y2,double x0,double y0){
    return (x1-x0)*(x2-x0)+(y1-y0)*(y2-y0);
}
//两点距离
double distance(point p1,point p2){
     return sqrt((p1.x-p2.x)*(p1.x-p2.x)+(p1.y-p2.y)*(p1.y-p2.y));
}
double distance(double x1,double y1,double x2,double y2){
    return sqrt((x1-x2)*(x1-x2)+(y1-y2)*(y1-y2));
}
//判三点共线
int dots inline(point p1, point p2, point p3){
    return zero(xmult(p1,p2,p3));
}
int dots inline(double x1,double y1,double x2,double y2,double x3,double y3){
    return zero(xmult(x1,y1,x2,y2,x3,y3));
}
//判点是否在线段上,包括端点
int dot_online_in(point p,line l){
    return zero(xmult(p,l.a,l.b))&&(l.a.x-p.x)*(l.b.x-p.x)\leqeps&&(l.a.y-p.y)*(l.b.y-p.y)\leqeps;
int dot online in(point p,point 11,point 12){
    return zero(xmult(p,11,12))&&(11.x-p.x)*(12.x-p.x)\leqeps&&(11.y-p.y)*(12.y-p.y)\leqeps;
int dot online in(double x,double y,double x1,double y1,double x2,double y2){
```

```
return zero(xmult(x,y,x1,y1,x2,y2))&&(x1-x)*(x2-x)<eps&&(y1-y)*(y2-y)<eps;
}
//判点是否在线段上,不包括端点
int dot online ex(point p,line l){
    return
dot online in(p,l)&&(!zero(p.x-l.a.x)||!zero(p.y-l.a.y))&&(!zero(p.x-l.b.x)||!zero(p.y-l.b.y));
int dot online ex(point p,point 11,point 12){
    return
dot online in(p,11,12)&&(!zero(p.x-11.x)||!zero(p.y-11.y))&&(!zero(p.x-12.x)||!zero(p.y-12.y));
int dot_online_ex(double x,double y,double x1,double y1,double x2,double y2){
dot\_online\_in(x,y,x1,y1,x2,y2)&&(!zero(x-x1)||!zero(y-y1))&&(!zero(x-x2)||!zero(y-y2));
//判两点在线段同侧,点在线段上返回0
int same side(point p1,point p2,line l){
    return xmult(l.a,p1,l.b)*xmult(l.a,p2,l.b)>eps;
int same_side(point p1,point p2,point l1,point l2){
    return xmult(11,p1,12)*xmult(11,p2,12)>eps;
}
//判两点在线段异侧,点在线段上返回0
int opposite_side(point p1,point p2,line l){
    return xmult(1.a,p1,1.b)*xmult(1.a,p2,1.b)<-eps;
int opposite side(point p1, point p2, point 11, point 12){
    return xmult(11,p1,12)*xmult(11,p2,12)<-eps;
}
//判两直线平行
int parallel(line u,line v){
    return zero((u.a.x-u.b.x)*(v.a.y-v.b.y)-(v.a.x-v.b.x)*(u.a.y-u.b.y));
int parallel(point u1,point u2,point v1,point v2){
    return zero((u1.x-u2.x)*(v1.y-v2.y)-(v1.x-v2.x)*(u1.y-u2.y));
}
//判两直线垂直
int perpendicular(line u,line v){
    return zero((u.a.x-u.b.x)*(v.a.x-v.b.x)+(u.a.y-u.b.y)*(v.a.y-v.b.y);
```

```
}
int perpendicular(point u1, point u2, point v1, point v2){
    return zero((u1.x-u2.x)*(v1.x-v2.x)+(u1.y-u2.y)*(v1.y-v2.y);
}
//判两线段相交,包括端点和部分重合
int intersect in(line u,line v){
    if (!dots inline(u.a,u.b,v.a)||!dots inline(u.a,u.b,v.b))
         return !same side(u.a,u.b,v)&&!same side(v.a,v.b,u);
    return dot online in(u.a,v)||dot online in(u.b,v)||dot online in(v.a,u)||dot online in(v.b,u);
}
int intersect in(point u1, point u2, point v1, point v2){
    if (!dots_inline(u1,u2,v1)||!dots_inline(u1,u2,v2))
         return !same side(u1,u2,v1,v2)&&!same side(v1,v2,u1,u2);
    return
dot online in(u1,v1,v2)||dot online in(u2,v1,v2)||dot online in(v1,u1,u2)||dot online in(v2,u1,u
2);
}
//判两线段相交,不包括端点和部分重合
int intersect ex(line u,line v){
    return opposite_side(u.a,u.b,v)&&opposite_side(v.a,v.b,u);
}
int intersect ex(point u1,point u2,point v1,point v2){
    return opposite side(u1,u2,v1,v2)&&opposite side(v1,v2,u1,u2);
}
//计算两直线交点,注意事先判断直线是否平行!
//线段交点请另外判线段相交(同时还是要判断是否平行!)
point intersection(line u,line v){
    point ret=u.a;
    double t=((u.a.x-v.a.x)*(v.a.y-v.b.y)-(u.a.y-v.a.y)*(v.a.x-v.b.x))
              /((u.a.x-u.b.x)*(v.a.y-v.b.y)-(u.a.y-u.b.y)*(v.a.x-v.b.x));
    ret.x+=(u.b.x-u.a.x)*t;
    ret.y+=(u.b.y-u.a.y)*t;
    return ret;
point intersection(point u1,point u2,point v1,point v2){
    point ret=u1;
    double t=((u1.x-v1.x)*(v1.y-v2.y)-(u1.y-v1.y)*(v1.x-v2.x))
              /((u1.x-u2.x)*(v1.y-v2.y)-(u1.y-u2.y)*(v1.x-v2.x));
    ret.x+=(u2.x-u1.x)*t;
    ret.y+=(u2.y-u1.y)*t;
    return ret;
```

```
}
//点到直线上的最近点
point ptoline(point p,line l){
     point t=p;
     t.x+=1.a.y-1.b.y,t.y+=1.b.x-1.a.x;
     return intersection(p,t,l.a,l.b);
}
point ptoline(point p,point 11,point 12){
     point t=p;
     t.x+=11.y-12.y,t.y+=12.x-11.x;
     return intersection(p,t,11,12);
}
//点到直线距离
double disptoline(point p,line l){
     return fabs(xmult(p,l.a,l.b))/distance(l.a,l.b);
}
double disptoline(point p,point 11,point 12){
     return fabs(xmult(p,l1,l2))/distance(l1,l2);
double disptoline(double x,double y,double x1,double y1,double x2,double y2){
     return fabs(xmult(x,y,x1,y1,x2,y2))/distance(x1,y1,x2,y2);
}
//点到线段上的最近点
point ptoseg(point p,line l){
     point t=p;
     t.x+=1.a.y-1.b.y,t.y+=1.b.x-1.a.x;
     if (xmult(l.a,t,p)*xmult(l.b,t,p)>eps)
          return distance(p,l.a)<distance(p,l.b)?l.a:l.b;
     return intersection(p,t,l.a,l.b);
point ptoseg(point p,point 11,point 12){
     point t=p;
     t.x+=11.y-12.y,t.y+=12.x-11.x;
     if (x mult(11,t,p)*x mult(12,t,p)>eps)
          return distance(p,l1)<distance(p,l2)?l1:l2;
     return intersection(p,t,11,12);
}
//点到线段距离
double disptoseg(point p,line l){
     point t=p;
```

```
t.x+=1.a.y-1.b.y,t.y+=1.b.x-1.a.x;
     if (xmult(l.a,t,p)*xmult(l.b,t,p)>eps)
          return distance(p,l.a)<distance(p,l.b)?distance(p,l.a):distance(p,l.b);
     return fabs(xmult(p,l.a,l.b))/distance(l.a,l.b);
}
double disptoseg(point p,point 11,point 12){
     point t=p;
     t.x+=11.y-12.y,t.y+=12.x-11.x;
     if (x mult(11,t,p)*x mult(12,t,p)>eps)
          return distance(p,11)<distance(p,12)?distance(p,11):distance(p,12);
     return fabs(xmult(p,11,12))/distance(11,12);
}
//矢量 V 以 P 为顶点逆时针旋转 angle 并放大 scale 倍
point rotate(point v,point p,double angle,double scale){
     point ret=p;
     v.x-=p.x,v.y-=p.y;
     p.x=scale*cos(angle);
     p.y=scale*sin(angle);
     ret.x+=v.x*p.x-v.y*p.y;
     ret.y+=v.x*p.y+v.y*p.x;
     return ret;
}
```

### 1.6 面积

```
#include <math.h>
struct point{double x,y;};
//计算 cross product (P1-P0)x(P2-P0)
double xmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
double xmult(double x1,double y1,double x2,double y2,double x0,double y0){
    return (x1-x0)*(y2-y0)-(x2-x0)*(y1-y0);
}
//计算三角形面积,输入三顶点
double area triangle(point p1,point p2,point p3){
    return fabs(xmult(p1,p2,p3))/2;
}
double area triangle(double x1,double y1,double x2,double y2,double x3,double y3){
    return fabs(xmult(x1,y1,x2,y2,x3,y3))/2;
}
```

```
//计算三角形面积,输入三边长
double area triangle(double a,double b,double c){
    double s=(a+b+c)/2;
    return sqrt(s*(s-a)*(s-b)*(s-c));
}
//计算多边形面积,顶点按顺时针或逆时针给出
double area polygon(int n,point* p){
    double s1=0,s2=0;
    int i:
    for (i=0; i< n; i++)
         s1+=p[(i+1)\%n].y*p[i].x,s2+=p[(i+1)\%n].y*p[(i+2)\%n].x;
    return fabs(s1-s2)/2;
}
       球面
1.7
#include <math.h>
const double pi=acos(-1);
//计算圆心角 lat 表示纬度,-90<=w<=90,lng 表示经度
//返回两点所在大圆劣弧对应圆心角,0<=angle<=pi
double angle(double lng1,double lat1,double lng2,double lat2){
    double dlng=fabs(lng1-lng2)*pi/180;
    while (dlng>=pi+pi)
         dlng-=pi+pi;
    if (dlng>pi)
         dlng=pi+pi-dlng;
    lat1*=pi/180, lat2*=pi/180;
    return acos(cos(lat1)*cos(lat2)*cos(dlng)+sin(lat1)*sin(lat2));
}
//计算距离,r 为球半径
double line dist(double r,double lng1,double lat1,double lng2,double lat2){
    double dlng=fabs(lng1-lng2)*pi/180;
    while (dlng>=pi+pi)
         dlng-=pi+pi;
    if (dlng>pi)
         dlng=pi+pi-dlng;
    lat1*=pi/180,lat2*=pi/180;
    return r*sqrt(2-2*(cos(lat1)*cos(lat2)*cos(dlng)+sin(lat1)*sin(lat2)));
```

```
//计算球面距离,r 为球半径
inline double sphere dist(double r,double lng1,double lat1,double lng2,double lat2){
     return r*angle(lng1,lat1,lng2,lat2);
}
1.8
        三角形
#include <math.h>
struct point{double x,y;};
struct line{point a,b;};
double distance(point p1,point p2){
     return sqrt((p1.x-p2.x)*(p1.x-p2.x)+(p1.y-p2.y)*(p1.y-p2.y));
}
point intersection(line u,line v){
     point ret=u.a;
     double t=((u.a.x-v.a.x)*(v.a.y-v.b.y)-(u.a.y-v.a.y)*(v.a.x-v.b.x))
               /((u.a.x-u.b.x)*(v.a.y-v.b.y)-(u.a.y-u.b.y)*(v.a.x-v.b.x));
     ret.x+=(u.b.x-u.a.x)*t;
     ret.y+=(u.b.y-u.a.y)*t;
     return ret;
}
//外心
point circumcenter(point a,point b,point c){
     line u,v;
     u.a.x = (a.x + b.x)/2;
     u.a.y=(a.y+b.y)/2;
     u.b.x=u.a.x-a.y+b.y;
     u.b.y=u.a.y+a.x-b.x;
     v.a.x = (a.x + c.x)/2;
     v.a.y=(a.y+c.y)/2;
     v.b.x=v.a.x-a.y+c.y;
     v.b.y=v.a.y+a.x-c.x;
     return intersection(u,v);
}
//内心
point incenter(point a,point b,point c){
```

line u,v; double m,n;

u.a=a;

m=atan2(b.y-a.y,b.x-a.x);

```
n=atan2(c.y-a.y,c.x-a.x);
    u.b.x=u.a.x+cos((m+n)/2);
    u.b.y=u.a.y+sin((m+n)/2);
    v.a=b;
    m=atan2(a.y-b.y,a.x-b.x);
    n=atan2(c.y-b.y,c.x-b.x);
    v.b.x=v.a.x+cos((m+n)/2);
    v.b.y=v.a.y+sin((m+n)/2);
    return intersection(u,v);
}
//垂心
point perpencenter(point a,point b,point c){
    line u,v;
    u.a=c;
    u.b.x=u.a.x-a.y+b.y;
    u.b.y=u.a.y+a.x-b.x;
    v.a=b;
    v.b.x=v.a.x-a.y+c.y;
    v.b.y=v.a.y+a.x-c.x;
    return intersection(u,v);
}
//重心
//到三角形三顶点距离的平方和最小的点
//三角形内到三边距离之积最大的点
point barycenter(point a,point b,point c){
    line u,v;
    u.a.x = (a.x + b.x)/2;
    u.a.y=(a.y+b.y)/2;
    u.b=c;
    v.a.x = (a.x + c.x)/2;
    v.a.y=(a.y+c.y)/2;
    v.b=b;
    return intersection(u,v);
}
//费马点
//到三角形三顶点距离之和最小的点
point fermentpoint(point a,point b,point c){
    point u,v;
    double step=fabs(a.x)+fabs(a.y)+fabs(b.x)+fabs(b.y)+fabs(c.x)+fabs(c.y);
    int i,j,k;
    u.x = (a.x + b.x + c.x)/3;
```

```
u.y=(a.y+b.y+c.y)/3;
while (step>1e-10)
    for (k=0;k<10;step/=2,k++)
        for (i=-1;i<=1;i++)
        for (j=-1;j<=1;j++){
            v.x=u.x+step*i;
            v.y=u.y+step*j;
            if
(distance(u,a)+distance(u,b)+distance(u,c)>distance(v,a)+distance(v,b)+distance(v,c))
            u=v;
        }
    return u;
}
```

### 1.9 三维几何

```
//三维几何函数库
#include <math.h>
#define eps 1e-8
#define zero(x) (((x)>0?(x):-(x))<eps)
struct point3{double x,y,z;};
struct line3{point3 a,b;};
struct plane3 {point3 a,b,c;};
//计算 cross product U x V
point3 xmult(point3 u,point3 v){
     point3 ret;
     ret.x=u.y*v.z-v.y*u.z;
     ret.y=u.z*v.x-u.x*v.z;
     ret.z=u.x*v.y-u.y*v.x;
     return ret;
}
//计算 dot product U.V
double dmult(point3 u,point3 v){
     return u.x*v.x+u.y*v.y+u.z*v.z;
}
//矢量差 U-V
point3 subt(point3 u,point3 v){
     point3 ret;
     ret.x=u.x-v.x;
     ret.y=u.y-v.y;
     ret.z=u.z-v.z;
```

```
return ret;
}
//取平面法向量
point3 pvec(plane3 s){
               return xmult(subt(s.a,s.b),subt(s.b,s.c));
point3 pvec(point3 s1,point3 s2,point3 s3){
               return xmult(subt(s1,s2),subt(s2,s3));
}
//两点距离,单参数取向量大小
double distance(point3 p1,point3 p2){
               return sqrt((p1.x-p2.x)*(p1.x-p2.x)+(p1.y-p2.y)*(p1.y-p2.y)+(p1.z-p2.z)*(p1.z-p2.z));
}
//向量大小
double vlen(point3 p){
               return sqrt(p.x*p.x+p.y*p.y+p.z*p.z);
}
//判三点共线
int dots inline(point3 p1,point3 p2,point3 p3){
               return vlen(xmult(subt(p1,p2),subt(p2,p3)))<eps;</pre>
}
//判四点共面
int dots onplane(point3 a,point3 b,point3 c,point3 d){
               return zero(dmult(pvec(a,b,c),subt(d,a)));
}
//判点是否在线段上,包括端点和共线
int dot_online_in(point3 p,line3 l){
               return\ zero(vlen(xmult(subt(p,l.a),subt(p,l.b))))\&\&(l.a.x-p.x)*(l.b.x-p.x) < eps\&\&(l.a.x-p.x) + (l.b.x-p.x) < eps\&\&(l.a.x-p.x) < eps\&\&(l.a.x-p
                               (l.a.y-p.y)*(l.b.y-p.y) < eps & & (l.a.z-p.z)*(l.b.z-p.z) < eps;
int dot online in(point3 p,point3 11,point3 12){
               return\ zero(vlen(xmult(subt(p,l1),subt(p,l2))))\&\&(l1.x-p.x)*(l2.x-p.x) < eps\&\&line(xmult(subt(p,l1),subt(p,l2))))\&\&(l1.x-p.x)*(l2.x-p.x) < eps\&\&line(xmult(subt(p,l1),subt(p,l2))))\&\&(l1.x-p.x)*(l2.x-p.x) < eps\&\&line(xmult(subt(p,l1),subt(p,l2))))\&\&(l1.x-p.x)*(l2.x-p.x) < eps\&\&line(xmult(subt(p,l1),subt(p,l2))))\&\&(l1.x-p.x)*(l2.x-p.x) < eps\&\&line(xmult(subt(p,l1),subt(p,l2))))\&\&(l1.x-p.x)*(l2.x-p.x) < eps\&\&line(xmult(subt(p,l1),subt(p,l2))))\&\&(l1.x-p.x) < eps\&\&line(xmult(subt(p,l2),subt(p,l2)))).
                               (11.y-p.y)*(12.y-p.y) < eps & & (11.z-p.z)*(12.z-p.z) < eps;
}
//判点是否在线段上,不包括端点
int dot_online_ex(point3 p,line3 l){
               return dot online in(p,l)&&(!zero(p.x-l.a.x)||!zero(p.y-l.a.y)||!zero(p.z-l.a.z))&&
```

```
(!zero(p.x-l.b.x)||!zero(p.y-l.b.y)||!zero(p.z-l.b.z));
int dot online ex(point3 p,point3 11,point3 12){
    return dot online in(p,11,12)&&(!zero(p,x-11.x)||!zero(p,y-11.y)||!zero(p,z-11.z))&&
         (!zero(p.x-12.x)||!zero(p.y-12.y)||!zero(p.z-12.z));
}
//判点是否在空间三角形上,包括边界,三点共线无意义
int dot inplane in(point3 p,plane3 s){
    return zero(vlen(xmult(subt(s.a,s.b),subt(s.a,s.c)))-vlen(xmult(subt(p,s.a),subt(p,s.b)))-
         vlen(xmult(subt(p,s.b),subt(p,s.c)))-vlen(xmult(subt(p,s.c),subt(p,s.a))));
int dot_inplane_in(point3 p,point3 s1,point3 s2,point3 s3){
    return zero(vlen(xmult(subt(s1,s2),subt(s1,s3)))-vlen(xmult(subt(p,s1),subt(p,s2)))-
         vlen(xmult(subt(p,s2),subt(p,s3)))-vlen(xmult(subt(p,s3),subt(p,s1))));
}
//判点是否在空间三角形上,不包括边界,三点共线无意义
int dot inplane ex(point3 p,plane3 s){
    return dot inplane in(p,s)&&vlen(xmult(subt(p,s.a),subt(p,s.b)))>eps&&
         vlen(xmult(subt(p,s.b),subt(p,s.c))) > eps & vlen(xmult(subt(p,s.c),subt(p,s.a))) > eps;
int dot inplane ex(point3 p,point3 s1,point3 s2,point3 s3){
    return dot inplane in(p,s1,s2,s3)&&vlen(xmult(subt(p,s1),subt(p,s2)))>eps&&
         vlen(xmult(subt(p,s2),subt(p,s3))) > eps&&vlen(xmult(subt(p,s3),subt(p,s1))) > eps;
}
//判两点在线段同侧,点在线段上返回 0,不共面无意义
int same_side(point3 p1,point3 p2,line3 l){
    return dmult(xmult(subt(1.a,1.b),subt(p1,1.b)),xmult(subt(1.a,1.b),subt(p2,1.b)))>eps;
}
int same side(point3 p1,point3 p2,point3 11,point3 12){
    return dmult(xmult(subt(11,12),subt(p1,12)),xmult(subt(11,12),subt(p2,12)))>eps;
}
//判两点在线段异侧,点在线段上返回 0,不共面无意义
int opposite side(point3 p1,point3 p2,line3 l){
    return dmult(xmult(subt(l.a,l.b),subt(p1,l.b)),xmult(subt(l.a,l.b),subt(p2,l.b)))<-eps;
int opposite side(point3 p1,point3 p2,point3 11,point3 12){
    return dmult(xmult(subt(11,12),subt(p1,12)),xmult(subt(11,12),subt(p2,12)))<-eps;
}
//判两点在平面同侧,点在平面上返回 0
```

```
int same side(point3 p1,point3 p2,plane3 s){
     return dmult(pvec(s),subt(p1,s.a))*dmult(pvec(s),subt(p2,s.a))>eps;
int same side(point3 p1,point3 p2,point3 s1,point3 s2,point3 s3){
    return dmult(pvec(s1,s2,s3),subt(p1,s1))*dmult(pvec(s1,s2,s3),subt(p2,s1))>eps;
}
//判两点在平面异侧,点在平面上返回 0
int opposite side(point3 p1,point3 p2,plane3 s){
     return dmult(pvec(s), subt(p1, s.a))*dmult(pvec(s), subt(p2, s.a)) <-eps;
}
int opposite side(point3 p1,point3 p2,point3 s1,point3 s2,point3 s3){
    return dmult(pvec(s1,s2,s3),subt(p1,s1))*dmult(pvec(s1,s2,s3),subt(p2,s1))<-eps;
}
//判两直线平行
int parallel(line3 u,line3 v){
     return vlen(xmult(subt(u.a,u.b),subt(v.a,v.b)))<eps;
int parallel(point3 u1,point3 u2,point3 v1,point3 v2){
     return vlen(xmult(subt(u1,u2),subt(v1,v2)))<eps;
}
//判两平面平行
int parallel(plane3 u,plane3 v){
     return vlen(xmult(pvec(u),pvec(v)))<eps;
int parallel(point3 u1,point3 u2,point3 u3,point3 v1,point3 v2,point3 v3){
    return vlen(xmult(pvec(u1,u2,u3),pvec(v1,v2,v3)))<eps;
}
//判直线与平面平行
int parallel(line3 l,plane3 s){
    return zero(dmult(subt(l.a,l.b),pvec(s)));
}
int parallel(point3 11,point3 12,point3 s1,point3 s2,point3 s3){
    return zero(dmult(subt(11,12),pvec(s1,s2,s3)));
}
//判两直线垂直
int perpendicular(line3 u,line3 v){
    return zero(dmult(subt(u.a,u.b),subt(v.a,v.b)));
int perpendicular(point3 u1,point3 u2,point3 v1,point3 v2){
```

```
return zero(dmult(subt(u1,u2),subt(v1,v2)));
}
//判两平面垂直
int perpendicular(plane3 u,plane3 v){
           return zero(dmult(pvec(u),pvec(v)));
int perpendicular(point3 u1,point3 u2,point3 u3,point3 v1,point3 v2,point3 v3){
           return zero(dmult(pvec(u1,u2,u3),pvec(v1,v2,v3)));
}
//判直线与平面平行
int perpendicular(line3 l,plane3 s){
           return vlen(xmult(subt(l.a,l.b),pvec(s)))<eps;
}
int perpendicular(point3 11, point3 12, point3 s1, point3 s2, point3 s3){
           return vlen(xmult(subt(11,12),pvec(s1,s2,s3)))<eps;
}
//判两线段相交,包括端点和部分重合
int intersect in(line3 u,line3 v){
           if (!dots_onplane(u.a,u.b,v.a,v.b))
                       return 0;
           if (!dots_inline(u.a,u.b,v.a)||!dots_inline(u.a,u.b,v.b))
                       return !same side(u.a,u.b,v)&&!same side(v.a,v.b,u);
           return dot online in(u.a,v)||dot online in(u.b,v)||dot online in(v.a,u)||dot online in(v.b,u);
}
int intersect in(point3 u1,point3 u2,point3 v1,point3 v2){
           if (!dots_onplane(u1,u2,v1,v2))
                       return 0;
           if (!dots inline(u1,u2,v1)||!dots inline(u1,u2,v2))
                       return !same side(u1,u2,v1,v2)&&!same side(v1,v2,u1,u2);
           return
dot_online_in(u1,v1,v2)||dot_online_in(u2,v1,v2)||dot_online_in(v1,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u
2);
}
//判两线段相交,不包括端点和部分重合
int intersect ex(line3 u,line3 v){
           return dots onplane(u.a,u.b,v.a,v.b)&&opposite side(u.a,u.b,v)&&opposite side(v.a,v.b,u);
}
int intersect ex(point3 u1,point3 u2,point3 v1,point3 v2){
dots onplane(u1,u2,v1,v2)&&opposite side(u1,u2,v1,v2)&&opposite side(v1,v2,u1,u2);
```

```
}
//判线段与空间三角形相交,包括交于边界和(部分)包含
int intersect in(line3 l,plane3 s){
    return !same side(l.a,l.b,s)&&!same side(s.a,s.b,l.a,l.b,s.c)&&
         !same side(s.b,s.c,l.a,l.b,s.a)&&!same side(s.c,s.a,l.a,l.b,s.b);
int intersect in(point3 11,point3 12,point3 s1,point3 s2,point3 s3){
    return !same side(11,12,s1,s2,s3)&&!same side(s1,s2,11,12,s3)&&
         !same side(s2,s3,11,12,s1)&&!same side(s3,s1,11,12,s2);
}
//判线段与空间三角形相交,不包括交于边界和(部分)包含
int intersect ex(line3 l,plane3 s){
    return opposite side(l.a,l.b,s)&&opposite side(s.a,s.b,l.a,l.b,s.c)&&
         opposite side(s.b,s.c,l.a,l.b,s.a)&&opposite side(s.c,s.a,l.a,l.b,s.b);
int intersect ex(point3 11,point3 12,point3 s1,point3 s2,point3 s3){
    return opposite side(11,12,s1,s2,s3)&&opposite side(s1,s2,11,12,s3)&&
         opposite side(s2,s3,l1,l2,s1)&&opposite side(s3,s1,l1,l2,s2);
}
//计算两直线交点,注意事先判断直线是否共面和平行!
//线段交点请另外判线段相交(同时还是要判断是否平行!)
point3 intersection(line3 u,line3 v){
    point3 ret=u.a;
    double t=((u.a.x-v.a.x)*(v.a.y-v.b.y)-(u.a.y-v.a.y)*(v.a.x-v.b.x))
              /((u.a.x-u.b.x)*(v.a.y-v.b.y)-(u.a.y-u.b.y)*(v.a.x-v.b.x));
    ret.x+=(u.b.x-u.a.x)*t;
    ret.y+=(u.b.y-u.a.y)*t;
    ret.z+=(u.b.z-u.a.z)*t;
    return ret;
point3 intersection(point3 u1,point3 u2,point3 v1,point3 v2){
    point3 ret=u1;
    double t=((u1.x-v1.x)*(v1.y-v2.y)-(u1.y-v1.y)*(v1.x-v2.x))
             /((u1.x-u2.x)*(v1.y-v2.y)-(u1.y-u2.y)*(v1.x-v2.x));
    ret.x += (u2.x - u1.x)*t;
    ret.y+=(u2.y-u1.y)*t;
    ret.z+=(u2.z-u1.z)*t;
    return ret;
}
```

//计算直线与平面交点,注意事先判断是否平行,并保证三点不共线!

```
//线段和空间三角形交点请另外判断
point3 intersection(line3 l,plane3 s){
    point3 ret=pvec(s);
    double t=(ret.x*(s.a.x-l.a.x)+ret.y*(s.a.y-l.a.y)+ret.z*(s.a.z-l.a.z))/
         (ret.x*(l.b.x-l.a.x)+ret.y*(l.b.y-l.a.y)+ret.z*(l.b.z-l.a.z));
    ret.x=l.a.x+(l.b.x-l.a.x)*t;
    ret.y=l.a.y+(l.b.y-l.a.y)*t;
    ret.z=l.a.z+(l.b.z-l.a.z)*t;
    return ret:
point3 intersection(point3 11, point3 12, point3 s1, point3 s2, point3 s3){
    point3 ret=pvec(s1,s2,s3);
    double t=(ret.x*(s1.x-l1.x)+ret.y*(s1.y-l1.y)+ret.z*(s1.z-l1.z))/
         (ret.x*(12.x-11.x)+ret.y*(12.y-11.y)+ret.z*(12.z-11.z));
    ret.x=11.x+(12.x-11.x)*t;
    ret.y=11.y+(12.y-11.y)*t;
    ret.z=11.z+(12.z-11.z)*t;
    return ret;
}
//计算两平面交线,注意事先判断是否平行,并保证三点不共线!
line3 intersection(plane3 u,plane3 v){
    line3 ret;
    c);
    ret.b=parallel(v.c,v.a,u.a,u.b,u.c)?intersection(v.b,v.c,u.a,u.b,u.c):intersection(v.c,v.a,u.a,u.b,u.
c);
    return ret;
line3 intersection(point3 u1,point3 u2,point3 u3,point3 v1,point3 v2,point3 v3){
    line3 ret;
    ret.a=parallel(v1,v2,u1,u2,u3)?intersection(v2,v3,u1,u2,u3):intersection(v1,v2,u1,u2,u3);
    ret.b=parallel(v3,v1,u1,u2,u3)?intersection(v2,v3,u1,u2,u3):intersection(v3,v1,u1,u2,u3);
    return ret;
}
//点到直线距离
double ptoline(point3 p,line3 l){
    return vlen(xmult(subt(p,l.a),subt(l.b,l.a)))/distance(l.a,l.b);
double ptoline(point3 p,point3 11,point3 12){
    return vlen(xmult(subt(p,11),subt(12,11)))/distance(11,12);
}
```

```
//点到平面距离
double ptoplane(point3 p,plane3 s){
    return fabs(dmult(pvec(s),subt(p,s.a)))/vlen(pvec(s));
}
double ptoplane(point3 p,point3 s1,point3 s2,point3 s3){
    return fabs(dmult(pvec(s1,s2,s3),subt(p,s1)))/vlen(pvec(s1,s2,s3));
}
//直线到直线距离
double linetoline(line3 u,line3 v){
    point3 n=xmult(subt(u.a,u.b),subt(v.a,v.b));
    return fabs(dmult(subt(u.a,v.a),n))/vlen(n);
double linetoline(point3 u1,point3 u2,point3 v1,point3 v2){
    point3 n=xmult(subt(u1,u2),subt(v1,v2));
    return fabs(dmult(subt(u1,v1),n))/vlen(n);
}
//两直线夹角 cos 值
double angle cos(line3 u,line3 v){
     return dmult(subt(u.a,u.b),subt(v.a,v.b))/vlen(subt(u.a,u.b))/vlen(subt(v.a,v.b));
}
double angle cos(point3 u1,point3 u2,point3 v1,point3 v2){
    return dmult(subt(u1,u2),subt(v1,v2))/vlen(subt(u1,u2))/vlen(subt(v1,v2));
}
//两平面夹角 cos 值
double angle cos(plane3 u,plane3 v){
    return dmult(pvec(u),pvec(v))/vlen(pvec(u))/vlen(pvec(v));
double angle cos(point3 u1,point3 u2,point3 u3,point3 v1,point3 v2,point3 v3){
    return dmult(pvec(u1,u2,u3),pvec(v1,v2,v3))/vlen(pvec(u1,u2,u3))/vlen(pvec(v1,v2,v3));
}
//直线平面夹角 sin 值
double angle sin(line3 l,plane3 s){
    return dmult(subt(l.a,l.b),pvec(s))/vlen(subt(l.a,l.b))/vlen(pvec(s));
}
double angle sin(point3 11,point3 12,point3 s1,point3 s2,point3 s3){
    return dmult(subt(11,12),pvec(s1,s2,s3))/vlen(subt(11,12))/vlen(pvec(s1,s2,s3));
}
```

# 1.10 凸包

```
#include <stdlib.h>
#define eps 1e-8
#define zero(x) (((x)>0?(x):-(x))<eps)
struct point{double x, y;};
//计算 cross product (P1-P0)x(P2-P0)
double xmult(point p1, point p2, point p0) {
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
//graham 算法顺时针构造包含所有共线点的凸包,O(nlogn)
point p1,p2;
int graham cp(const void* a,const void* b){
    double ret=xmult(*((point*)a),*((point*)b),p1);
    return zero(ret)?(xmult(*((point*)a),*((point*)b),p2)>0?1:-1):(ret>0?1:-1);
}
void graham(int n,point* p,int& s,point* ch){
    int i,k=0;
    for (p1=p2=p[0],i=1;i < n;p2.x+=p[i].x,p2.y+=p[i].y,i++)
        if(p1.y-p[i].y>eps||(zero(p1.y-p[i].y)&&p1.x>p[i].x))
             p1=p[k=i];
    p2.x/=n, p2.y/=n;
    p[k]=p[0],p[0]=p1;
    qsort(p+1,n-1,sizeof(point),graham cp);
    for (ch[0]=p[0],ch[1]=p[1],ch[2]=p[2],s=i=3;i < n;ch[s++]=p[i++])
        for (s>2\&xmult(ch[s-2],p[i],ch[s-1])<-eps;s--);
}
//构造凸包接口函数,传入原始点集大小 n,点集 p(p 原有顺序被打乱!)
//返回凸包大小,凸包的点在 convex 中
//参数 maxsize 为 1 包含共线点,为 0 不包含共线点,缺省为 1
//参数 clockwise 为 1 顺时针构造,为 0 逆时针构造,缺省为 1
//在输入仅有若干共线点时算法不稳定,可能有此类情况请另行处理!
//不能去掉点集中重合的点
int graham(int n,point* p,point* convex,int maxsize=1,int dir=1){
    point* temp=new point[n];
    int s,i;
    graham(n,p,s,temp);
    for (convex[0]=temp[0],n=1,i=(dir?1:(s-1));dir?(i<s):i;i+=(dir?1:-1))
        if (maxsize||!zero(xmult(temp[i-1],temp[i],temp[(i+1)%s])))
             convex[n++]=temp[i];
    delete []temp;
    return n;
```

### 1.11 网格

```
#define abs(x) ((x)>0?(x):-(x))
struct point{int x,y;};
int gcd(int a,int b){
     return b?gcd(b,a%b):a;
}
//多边形上的网格点个数
int grid onedge(int n,point* p){
     int i,ret=0;
     for (i=0;i< n;i++)
         ret+=gcd(abs(p[i].x-p[(i+1)\%n].x),abs(p[i].y-p[(i+1)\%n].y));
     return ret;
}
//多边形内的网格点个数
int grid_inside(int n,point* p){
     int i,ret=0;
     for (i=0;i<n;i++)
         ret+=p[(i+1)\%n].y*(p[i].x-p[(i+2)\%n].x);
     return (abs(ret)-grid onedge(n,p))/2+1;
}
```

#### 1.12 圆

```
#include <math.h>
#define eps 1e-8
struct point{double x,y;};

double xmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}

double distance(point p1,point p2){
    return sqrt((p1.x-p2.x)*(p1.x-p2.x)+(p1.y-p2.y)*(p1.y-p2.y));
}

double disptoline(point p,point l1,point l2){
    return fabs(xmult(p,l1,l2))/distance(l1,l2);
}
```

```
point intersection(point u1, point u2, point v1, point v2){
    point ret=u1;
    double t=((u1.x-v1.x)*(v1.y-v2.y)-(u1.y-v1.y)*(v1.x-v2.x))
             /((u1.x-u2.x)*(v1.y-v2.y)-(u1.y-u2.y)*(v1.x-v2.x));
    ret.x += (u2.x - u1.x) *t;
    ret.y+=(u2.y-u1.y)*t;
    return ret;
}
//判直线和圆相交,包括相切
int intersect line circle(point c,double r,point 11,point 12){
    return disptoline(c,11,12)<r+eps;
}
//判线段和圆相交,包括端点和相切
int intersect seg circle(point c,double r,point 11,point 12){
    double t1=distance(c,11)-r,t2=distance(c,12)-r;
    point t=c;
    if (t1 \le eps || t2 \le eps)
         return t1>-eps||t2>-eps;
    t.x+=11.y-12.y;
    t.y+=12.x-11.x;
    return xmult(11,c,t)*xmult(12,c,t)<eps&&disptoline(c,11,12)-r<eps;
}
//判圆和圆相交,包括相切
int intersect_circle_circle(point c1,double r1,point c2,double r2){
    return distance(c1,c2)<r1+r2+eps&&distance(c1,c2)>fabs(r1-r2)-eps;
}
//计算圆上到点 p 最近点,如 p 与圆心重合,返回 p 本身
point dot to circle(point c,double r,point p){
    point u,v;
    if (distance(p,c)<eps)
         return p;
    u.x=c.x+r*fabs(c.x-p.x)/distance(c,p);
    u.y=c.y+r*fabs(c.y-p.y)/distance(c,p)*((c.x-p.x)*(c.y-p.y)<0?-1:1);
    v.x=c.x-r*fabs(c.x-p.x)/distance(c,p);
    v.y=c.y-r*fabs(c.y-p.y)/distance(c,p)*((c.x-p.x)*(c.y-p.y)<0?-1:1);
    return distance(u,p)<distance(v,p)?u:v;
}
//计算直线与圆的交点,保证直线与圆有交点
//计算线段与圆的交点可用这个函数后判点是否在线段上
```

```
void intersection line circle(point c,double r,point 11,point 12,point& p1,point& p2){
    point p=c;
    double t;
    p.x+=11.y-12.y;
    p.y+=12.x-11.x;
    p=intersection(p,c,l1,l2);
    t=sqrt(r*r-distance(p,c)*distance(p,c))/distance(11,12);
    p1.x=p.x+(12.x-11.x)*t;
    p1.y=p.y+(12.y-11.y)*t;
    p2.x=p.x-(12.x-11.x)*t;
    p2.y=p.y-(12.y-11.y)*t;
}
//计算圆与圆的交点、保证圆与圆有交点、圆心不重合
void intersection circle circle(point c1,double r1,point c2,double r2,point& p1,point& p2){
    point u,v;
    double t;
    t=(1+(r_1*r_1-r_2*r_2)/distance(c_1,c_2)/distance(c_1,c_2))/2;
    u.x=c1.x+(c2.x-c1.x)*t;
    u.y=c1.y+(c2.y-c1.y)*t;
    v.x=u.x+c1.y-c2.y;
    v.y=u.y-c1.x+c2.x;
    intersection line circle(c1,r1,u,v,p1,p2);
}
```

#### 1.13 整数函数

```
//整数几何函数库
//注意某些情况下整数运算会出界!
#define sign(a) ((a)>0?1:(((a)<0?-1:0)))
struct point {int x,y;};
struct line{point a,b;};

//计算 cross product (P1-P0)x(P2-P0)
int xmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
int xmult(int x1,int y1,int x2,int y2,int x0,int y0){
    return (x1-x0)*(y2-y0)-(x2-x0)*(y1-y0);
}

//计算 dot product (P1-P0).(P2-P0)
int dmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.x-p0.x)+(p1.y-p0.y)*(p2.y-p0.y);
```

```
}
int dmult(int x1,int y1,int x2,int y2,int x0,int y0){
    return (x1-x0)*(x2-x0)+(y1-y0)*(y2-y0);
}
//判三点共线
int dots inline(point p1, point p2, point p3){
    return !xmult(p1,p2,p3);
}
int dots inline(int x1,int y1,int x2,int y2,int x3,int y3){
    return !xmult(x1,y1,x2,y2,x3,y3);
}
//判点是否在线段上,包括端点和部分重合
int dot online in(point p,line l){
    return !xmult(p,l.a,l.b)&&(l.a.x-p.x)*(l.b.x-p.x)\leq 0&&(l.a.y-p.y)*(l.b.y-p.y)\leq 0;
int dot online in(point p,point 11,point 12){
    return !xmult(p,11,12)&&(11.x-p.x)*(12.x-p.x) \le 0&&(11.y-p.y)*(12.y-p.y) \le 0;
int dot online in(int x,int y,int x1,int y1,int x2,int y2){
    return !xmult(x,y,x1,y1,x2,y2)&&(x1-x)*(x2-x)<=0&&(y1-y)*(y2-y)<=0;
}
//判点是否在线段上,不包括端点
int dot online ex(point p,line l){
     return dot_online_in(p,l)&&(p.x!=l.a.x||p.y!=l.a.y)&&(p.x!=l.b.x||p.y!=l.b.y);
}
int dot_online_ex(point p,point 11,point 12){
     return dot online in(p,11,12)&&(p.x!=11.x||p.y!=11.y)&&(p.x!=12.x||p.y!=12.y);
}
int dot online ex(int x,int y,int x1,int y1,int x2,int y2){
     return dot_online_in(x,y,x1,y1,x2,y2)&&(x!=x1||y!=y1)&&(x!=x2||y!=y2);
}
//判两点在直线同侧,点在直线上返回 0
int same side(point p1,point p2,line l){
    return sign(xmult(l.a,p1,l.b))*xmult(l.a,p2,l.b)>0;
int same side(point p1,point p2,point 11,point 12){
     return sign(xmult(11,p1,12))*xmult(11,p2,12)>0;
}
//判两点在直线异侧,点在直线上返回 0
```

```
int opposite side(point p1,point p2,line l){
            return sign(xmult(l.a,p1,l.b))*xmult(l.a,p2,l.b)<0;
int opposite side(point p1, point p2, point 11, point 12) {
           return sign(xmult(11,p1,12))*xmult(11,p2,12)<0;
}
//判两直线平行
int parallel(line u,line v){
           return (u.a.x-u.b.x)*(v.a.y-v.b.y)==(v.a.x-v.b.x)*(u.a.y-u.b.y);
int parallel(point u1, point u2, point v1, point v2){
           return (u1.x-u2.x)*(v1.y-v2.y)==(v1.x-v2.x)*(u1.y-u2.y);
}
//判两直线垂直
int perpendicular(line u,line v){
            return (u.a.x-u.b.x)*(v.a.x-v.b.x)==-(u.a.y-u.b.y)*(v.a.y-v.b.y);
int perpendicular(point u1, point u2, point v1, point v2){
            return (u1.x-u2.x)*(v1.x-v2.x)==-(u1.y-u2.y)*(v1.y-v2.y);
}
//判两线段相交,包括端点和部分重合
int intersect in(line u,line v){
           if (!dots inline(u.a,u.b,v.a)||!dots inline(u.a,u.b,v.b))
                       return !same_side(u.a,u.b,v)&&!same_side(v.a,v.b,u);
           return dot online in(u.a,v)||dot online in(u.b,v)||dot online in(v.a,u)||dot online in(v.b,u);
}
int intersect in(point u1, point u2, point v1, point v2){
           if (!dots inline(u1,u2,v1)||!dots inline(u1,u2,v2))
                       return !same side(u1,u2,v1,v2)&&!same side(v1,v2,u1,u2);
           return
dot_online_in(u1,v1,v2)||dot_online_in(u2,v1,v2)||dot_online_in(v1,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u2)||dot_online_in(v2,u1,u
2);
}
//判两线段相交,不包括端点和部分重合
int intersect ex(line u,line v){
            return opposite side(u.a,u.b,v)&&opposite side(v.a,v.b,u);
}
int intersect ex(point u1,point u2,point v1,point v2){
           return opposite_side(u1,u2,v1,v2)&&opposite_side(v1,v2,u1,u2);
```

#### 2、组合

#### 2.1 组合公式

```
1. C(m,n)=C(m,m-n)
2. C(m,n)=C(m-1,n)+C(m-1,n-1)
derangement D(n) = n!(1 - 1/1! + 1/2! - 1/3! + ... + (-1)^n/n!)
                    = (n-1)(D(n-2) - D(n-1))
              Q(n) = D(n) + D(n-1)
求和公式,k=1..n
1. sum( k ) = n(n+1)/2
2. sum( 2k-1 ) = n^2
3. sum( k^2 ) = n(n+1)(2n+1)/6
4. sum( (2k-1)^2 ) = n(4n^2-1)/3
5. sum( k^3 ) = (n(n+1)/2)^2
6. sum( (2k-1)^3 ) = n^2(2n^2-1)
7. sum( k^4) = n(n+1)(2n+1)(3n^2+3n-1)/30
8. sum( k^5 ) = n^2(n+1)^2(2n^2+2n-1)/12
9. sum( k(k+1) ) = n(n+1)(n+2)/3
10. sum( k(k+1)(k+2) ) = n(n+1)(n+2)(n+3)/4
12. sum( k(k+1)(k+2)(k+3) ) = n(n+1)(n+2)(n+3)(n+4)/5
```

### 2.2 排列组合生成

```
//gen_perm 产生字典序排列 P(n,m)
//gen_comb 产生字典序组合 C(n,m)
//gen_perm_swap 产生相邻位对换全排列 P(n,n)
//产生元素用 1..n 表示
//dummy 为产生后调用的函数,传入 a[]和 n,a[0]..a[n-1]为一次产生的结果
#define MAXN 100
int count;

#include <iostream.h>
void dummy(int* a,int n){
    int i;
    cout<<count++<<":";
    for (i=0;i<n-1;i++)
        cout<<a[i]<<'';
    cout<<a[n-1]<<endl;
}
```

```
void _gen_perm(int* a,int n,int m,int l,int* temp,int* tag){
     int i;
    if(l==m)
         dummy(temp,m);
     else
         for (i=0;i<n;i++)
              if (!tag[i]) {
                   temp[1]=a[i],tag[i]=1;
                   _gen_perm(a,n,m,l+1,temp,tag);
                   tag[i]=0;
              }
}
void gen_perm(int n,int m){
     int a[MAXN],temp[MAXN],tag[MAXN]={0},i;
     for (i=0;i<n;i++)
         a[i]=i+1;
     _gen_perm(a,n,m,0,temp,tag);
}
void _gen_comb(int* a,int s,int e,int m,int& count,int* temp){
     int i;
     if (!m)
         dummy(temp,count);
     else
         for (i=s; i \le e-m+1; i++){
              temp[count++]=a[i];
               _gen_comb(a,i+1,e,m-1,count,temp);
              count--;
         }
}
void gen_comb(int n,int m){
     int a[MAXN],temp[MAXN],count=0,i;
     for (i=0;i<n;i++)
         a[i]=i+1;
     _gen_comb(a,0,n-1,m,count,temp);
}
void _gen_perm_swap(int* a,int n,int l,int* pos,int* dir){
     int i,p1,p2,t;
     if(l==n)
         dummy(a,n);
```

```
else{
         _gen_perm_swap(a,n,l+1,pos,dir);
         for (i=0;i<1;i++)
              p2=(p1=pos[1])+dir[1];
              t=a[p1],a[p1]=a[p2],a[p2]=t;
              pos[a[p1]-1]=p1,pos[a[p2]-1]=p2;
              gen_perm_swap(a,n,l+1,pos,dir);
         }
         dir[1]=-dir[1];
    }
}
void gen_perm_swap(int n){
    int a[MAXN],pos[MAXN],dir[MAXN],i;
    for (i=0;i<n;i++)
         a[i]=i+1,pos[i]=i,dir[i]=-1;
     gen_perm_swap(a,n,0,pos,dir);
}
```

# 2.3 生成 gray 码

```
//生成 reflected gray code
//每次调用 gray 取得下一个码
//000...000 是第一个码,100...000 是最后一个码
void gray(int n,int *code) {
    int t=0,i;
    for (i=0;i<n;t+=code[i++]);
    if (t&1)
        for (n--;!code[n];n--);
    code[n-1]=1-code[n-1];
}
```

# 2.4 置换(polya)

```
//求置换的循环节,polya 原理
//perm[0..n-1]为 0..n-1 的一个置换(排列)
//返回置换最小周期,num 返回循环节个数
#define MAXN 1000

int gcd(int a,int b){
    return b?gcd(b,a%b):a;
}
```

```
int polya(int* perm,int n,int& num) {
    int i,j,p,v[MAXN]={0},ret=1;
    for (num=i=0;i<n;i++)
        if (!v[i]) {
        for (num++,j=0,p=i;!v[p=perm[p]];j++)
            v[p]=1;
        ret*=j/gcd(ret,j);
        }
    return ret;
}</pre>
```

#### 2.5 字典序全排列

```
//字典序全排列与序号的转换
int perm2num(int n,int *p){
     int i,j,ret=0,k=1;
     for (i=n-2;i>=0;k*=n-(i--))
          for (j=i+1;j< n;j++)
               if(p[j] < p[i])
                    ret+=k;
     return ret;
}
void num2perm(int n,int *p,int t){
     int i,j;
     for (i=n-1;i>=0;i--)
          p[i]=t\%(n-i),t/=n-i;
     for (i=n-1;i;i--)
          for (j=i-1;j>=0;j--)
               if (p[i] \le p[i])
                    p[i]++;
}
```

## 2.6 字典序组合

```
//字典序组合与序号的转换
//comb 为组合数 C(n,m),必要时换成大数,注意处理 C(n,m)=0|n<m
int comb(int n,int m){
    int ret=1,i;
    m=m<(n-m)?m:(n-m);
    for (i=n-m+1;i<=n;ret*=(i++));
    for (i=1;i<=m;ret/=(i++));
    return m<0?0:ret;
```

```
}
int comb2num(int n,int m,int *c){
    int ret=comb(n,m),i;
    for (i=0;i<m;i++)
        ret-=comb(n-c[i],m-i);
    return ret;
}

void num2comb(int n,int m,int* c,int t){
    int i,j=1,k;
    for (i=0;i<m;c[i++]=j++)
        for (;t>(k=comb(n-j,m-i-1));t-=k,j++);
}
```

#### 3、结构

## 3.1 并查集

```
//带路径压缩的并查集,用于动态维护查询等价类
//图论算法中动态判点集连通常用
//维护和查询复杂度略大于 O(1)
//集合元素取值 1..MAXN-1(注意 0 不能用!),默认不等价
#include <string.h>
#define MAXN 100000
#define \_ufind\_run(x) for(;p[t=x];x=p[x],p[t]=(p[x]?p[x]:x))
#define run both ufind run(i); ufind run(j)
struct ufind{
    int p[MAXN],t;
    void init(){memset(p,0,sizeof(p));}
    void set friend(int i,int j){ run both;p[i]=(i==j?0:j);}
    int is_friend(int i,int j){_run_both;return i==j&&i;}
};
//带路径压缩的并查集扩展形式
//用于动态维护查询 friend-enemy 型等价类
//维护和查询复杂度略大于 O(1)
//集合元素取值 1..MAXN-1(注意 0 不能用!),默认无关
#include <string.h>
#define MAXN 100000
#define sig(x) ((x)>0?1:-1)
```

```
#define abs(x) ((x)>0?(x):-(x))
#define
                                                                           ufind run(x)
for(p[t=abs(x)]; x=sig(x)*p[abs(x)], p[t]=sig(p[t])*(p[abs(x)]; p[abs(x)]; abs(p[t])))
#define run both ufind run(i); ufind run(j)
#define set side(x) p[abs(i)]=sig(i)*(abs(i)==abs(j)?0:(x)*j)
#define judge side(x) (i==(x)*j\&\&i)
struct ufind{
    int p[MAXN],t;
    void init(){memset(p,0,sizeof(p));}
    int set friend(int i,int j){ run both; set side(1);return! judge side(-1);}
    int set enemy(int i,int j){ run both; set side(-1);return! judge side(1);}
    int is_friend(int i,int j){ run_both;return_judge_side(1);}
    int is enemy(int i,int j){ run both;return judge side(-1);}
};
3.2 堆
//二分堆(binary)
//可插入,获取并删除最小(最大)元素,复杂度均 O(logn)
//可更改元素类型,修改比较符号或换成比较函数
#define MAXN 10000
#define cp(a,b)((a)<(b))
typedef int elem t;
struct heap{
    elem_t h[MAXN];
    int n,p,c;
    void init(){n=0;}
    void ins(elem t e){
        for (p=++n;p>1&& cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
        h[p]=e;
    }
    int del(elem t& e){
        if (!n) return 0;
        for
(e=h[p=1],c=2;c<n\&\&\_cp(h[c+=(c< n-1\&\&\_cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
        h[p]=h[n--];return 1;
    }
};
//映射二分堆(mapped)
//可插入,获取并删除任意元素,复杂度均 O(logn)
//插入时提供一个索引值.删除时按该索引删除.获取并删除最小元素时一起获得该索引
```

```
//索引值范围 0..MAXN-1,不能重复,不负责维护索引的唯一性,不在此返回请另外映射
//主要用于图论算法,该索引值可以是节点的下标
//可更改元素类型,修改比较符号或换成比较函数
#define MAXN 10000
#define cp(a,b)((a)<(b))
typedef int elem t;
struct heap{
    elem t h[MAXN];
    int ind[MAXN],map[MAXN],n,p,c;
    void init(){n=0;}
    void ins(int i,elem t e){
        for (p=++n;p>1\&\&\_cp(e,h[p>>1]);h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
        h[map[ind[p]=i]=p]=e;
    }
    int del(int i,elem t& e){
        i=map[i];if(i<1||i>n) return 0;
        for (e=h[p=i];p>1;h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
        for
(c=2;c<n\&\&\_cp(h[c+=(c< n-1\&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<< n\&\&\_cp(h[c+1],h[c]))
=1);
        h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
    }
    int delmin(int& i,elem t& e){
        if (n<1) return 0; i=ind[1];
(e=h[p=1],c=2;c<n\&\&\_cp(h[c+=(c<n-1\&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c]
],p=c,c<<=1);
        h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
};
```

#### 3.3 线段树

线段树应用:

求面积:

- 1) 坐标离散化
- 2) 垂直边按 x 坐标排序
- 3) 从左往右用线段树处理垂直边 累计每个离散 x 区间长度和线段树长度的乘积 求周长:
- 1) 坐标离散化
- 2) 垂直边按 x 坐标排序, 第二关键字为入边优于出边
- 3) 从左往右用线段树处理垂直边

在每个离散点上先加入所有入边,累计线段树长度变化值再删除所有出边,累计线段树长度变化值

- 4) 水平边按 y 坐标排序, 第二关键字为入边优于出边
- 5) 从上往下用线段树处理水平边 在每个离散点上先加入所有入边,累计线段树长度变化值 再删除所有出边,累计线段树长度变化值

```
//线段树
//可以处理加入边和删除边不同的情况
//inc seg 和 dec seg 用于加入边
//seg len 求长度
//t 传根节点(一律为 1)
//10,r0 传树的节点范围(一律为 1..t)
//l,r 传线段(端点)
#define MAXN 10000
struct segtree{
     int n,cnt[MAXN],len[MAXN];
     segtree(int t):n(t){
         for (int i=1; i <=t; i++)
              cnt[i]=len[i]=0;
     };
     void update(int t,int l,int r);
     void inc seg(int t,int 10,int r0,int 1,int r);
     void dec seg(int t,int l0,int r0,int l,int r);
     int seg len(int t,int 10,int r0,int 1,int r);
};
int length(int l,int r){
     return r-1;
}
void segtree::update(int t,int l,int r){
     if (cnt[t]||r-l==1)
         len[t]=length(l,r);
     else
         len[t]=len[t+t]+len[t+t+1];
}
void segtree::inc seg(int t,int 10,int r0,int 1,int r){
     if (10==1\&\&r0==r)
         cnt[t]++;
     else{
         int m0=(10+r0)>>1;
         if (1<m0)
```

```
inc seg(t+t,10,m0,1,m0 < r?m0:r);
          if (r>m0)
               inc seg(t+t+1,m0,r0,m0>1?m0:l,r);
          if(cnt[t+t]&&cnt[t+t+1]){
               cnt[t+t]--;
               update(t+t,l0,m0);
               cnt[t+t+1]--;
               update(t+t+1,m0,r0);
               cnt[t]++;
     }
     update(t,l0,r0);
}
void segtree::dec_seg(int t,int l0,int r0,int l,int r){
     if (10==1\&\&r0==r)
          cnt[t]--;
     else if (cnt[t]){
          cnt[t]--;
          if (l>l0)
               inc seg(t,10,r0,10,1);
          if (r<r0)
               inc_seg(t,10,r0,r,r0);
     }
     else{
          int m0=(10+r0)>>1;
          if(1 \le m0)
               dec_seg(t+t,l0,m0,l,m0<r?m0:r);
          if (r>m0)
               dec seg(t+t+1,m0,r0,m0>1?m0:l,r);
     }
     update(t,10,r0);
}
int segtree::seg len(int t,int l0,int r0,int l,int r){
     if(cnt[t]||(10==1\&\&r0==r))
          return len[t];
     else{
          int m0=(10+r0)>>1,ret=0;
          if(1 \le m0)
               ret+=seg_len(t+t,l0,m0,l,m0 < r?m0:r);
          if (r > m0)
               ret+=seg_len(t+t+1,m0,r0,m0>1?m0:l,r);
          return ret;
```

```
}
}
//线段树扩展
//可以计算长度和线段数
//可以处理加入边和删除边不同的情况
//inc seg 和 dec seg 用于加入边
//seg_len 求长度,seg_cut 求线段数
//t 传根节点(一律为 1)
//l0,r0 传树的节点范围(一律为 1..t)
//l,r 传线段(端点)
#define MAXN 10000
struct segtree{
     int n,cnt[MAXN],len[MAXN],cut[MAXN],bl[MAXN],br[MAXN];
     segtree(int t):n(t){
          for (int i=1; i <=t; i++)
               cnt[i]=len[i]=cut[i]=bl[i]=br[i]=0;
     };
     void update(int t,int l,int r);
     void inc_seg(int t,int 10,int r0,int 1,int r);
     void dec seg(int t,int l0,int r0,int l,int r);
     int seg_len(int t,int l0,int r0,int l,int r);
     int seg cut(int t,int 10,int r0,int 1,int r);
};
int length(int l,int r){
     return r-1;
}
void segtree::update(int t,int l,int r){
     if (cnt[t]||r-l==1)
          len[t] = length(l,r), cut[t] = bl[t] = br[t] = 1;
     else{
          len[t]=len[t+t]+len[t+t+1];
          \operatorname{cut}[t] = \operatorname{cut}[t+t] + \operatorname{cut}[t+t+1];
          if (br[t+t]&&bl[t+t+1])
               cut[t]--;
          bl[t]=bl[t+t],br[t]=br[t+t+1];
     }
}
void segtree::inc_seg(int t,int 10,int r0,int l,int r){
     if (10==1\&\&r0==r)
          cnt[t]++;
```

```
else{
          int m0=(10+r0)>>1;
          if (1<m0)
               inc seg(t+t,10,m0,1,m0 < r?m0:r);
          if (r>m0)
               inc seg(t+t+1,m0,r0,m0>1?m0:1,r);
          if(cnt[t+t]&&cnt[t+t+1]){
               cnt[t+t]--;
               update(t+t,l0,m0);
               cnt[t+t+1]--;
               update(t+t+1,m0,r0);
               cnt[t]++;
          }
     }
     update(t,l0,r0);
}
void segtree::dec_seg(int t,int l0,int r0,int l,int r){
     if (10==1\&\&r0==r)
          cnt[t]--;
     else if (cnt[t]){
          cnt[t]--;
          if (l>l0)
               inc_seg(t,10,r0,10,1);
          if (r<r0)
               inc seg(t,10,r0,r,r0);
     }
     else{
          int m0=(10+r0)>>1;
          if (1<m0)
               dec_seg(t+t,l0,m0,l,m0<r?m0:r);
          if (r>m0)
               dec_seg(t+t+1,m0,r0,m0>1?m0:l,r);
     }
     update(t,10,r0);
}
int segtree::seg_len(int t,int l0,int r0,int l,int r){
     if (cnt[t]||(10==1\&\&r0==r))
          return len[t];
     else{
          int m0=(10+r0)>>1,ret=0;
          if(1 \le m0)
               ret+=seg_len(t+t,l0,m0,l,m0 < r?m0:r);
```

```
if(r>m0)
               ret = seg len(t+t+1,m0,r0,m0>1?m0:1,r);
          return ret;
     }
}
int segtree::seg cut(int t,int 10,int r0,int 1,int r){
     if (cnt[t])
          return 1;
     if (10==1\&\&r0==r)
          return cut[t];
     else{
          int m0=(10+r0)>>1,ret=0;
          if(1 \le m0)
               ret+=seg_cut(t+t,l0,m0,l,m0<r?m0:r);
          if (r > m0)
               ret+=seg cut(t+t+1,m0,r0,m0>1?m0:1,r);
          if(1 < m0 & r > m0 & br[t+t] & bl[t+t+1])
               ret--;
          return ret;
}
```

#### 3.4 子段和

```
//求 sum{[0..n-1]}
//维护和查询复杂度均为 O(logn)
//用于动态求子段和,数组内容保存在 sum.a[]中
//可以改成其他数据类型
#include <string.h>
#define lowbit(x) ((x)&((x)^((x)-1)))
#define MAXN 10000
typedef int elem_t;

struct sum{
    elem_t a[MAXN],c[MAXN],ret;
    int n;
    void init(int i){memset(a,0,sizeof(a));memset(c,0,sizeof(c));n=i;}
    void update(int i,elem_t v){for (v-=a[i],a[i++]+=v;i<=n;c[i-1]+=v,i+=lowbit(i));}
    elem_t query(int i){for (ret=0;i;ret+=c[i-1],i^=lowbit(i));return ret;}
};
```

### 3.5 子阵和

```
//求 sum{a[0..m-1][0..n-1]}
//维护和查询复杂度均为 O(logm*logn)
//用于动态求子阵和,数组内容保存在 sum.a[][]中
//可以改成其他数据类型
#include <string.h>
#define lowbit(x) ((x)&((x)^{(x)-1)})
#define MAXN 100
typedef int elem t;
struct sum{
    elem t a[MAXN][MAXN],c[MAXN][MAXN],ret;
    int m,n,t;
    void init(int i,int j){memset(a,0,sizeof(a));memset(c,0,sizeof(c));m=i,n=j;}
    void update(int i,int j,elem_t v){
         for (v-=a[i][j],a[i++][j++]+=v,t=j;i \le m;i+=lowbit(i))
              for (j=t; j \le n; c[i-1][j-1] + = v, j + = lowbit(j));
    }
    elem t query(int i,int j){
         for (ret=0,t=j;i;i^=lowbit(i))
              for (j=t;j;ret+=c[i-1][j-1],j^=lowbit(j));
         return ret;
    }
};
```

#### 4、数论

#### 4.1 阶乘最后非 0 位

```
//求阶乘最后非零位,复杂度 O(nlogn)
//返回该位,n 以字符串方式传入
#include <string.h>
#define MAXN 10000

int lastdigit(char* buf) {
    const int mod[20]={1,1,2,6,4,2,2,4,2,8,4,4,8,4,6,8,8,6,8,2};
    int len=strlen(buf),a[MAXN],i,c,ret=1;
    if (len==1)
        return mod[buf[0]-'0'];
    for (i=0;i<len;i++)
        a[i]=buf[len-1-i]-'0';
    for (;len;len-=!a[len-1]) {
        ret=ret*mod[a[1]%2*10+a[0]]%5;
        for (c=0,i=len-1;i>=0;i--)
```

```
c = c*10 + a[i], a[i] = c/5, c\% = 5; \} return \ ret + ret\% 2*5; \}
```

### 4.2 模线性方程组

```
#ifdef WIN32
typedef int64 i64;
#else
typedef long long i64;
#endif
//扩展 Euclid 求解 gcd(a,b)=ax+by
int ext gcd(int a,int b,int& x,int& y){
    int t,ret;
    if (!b) {
         x=1,y=0;
         return a;
    }
    ret=ext_gcd(b,a%b,x,y);
    t=x,x=y,y=t-a/b*y;
    return ret;
}
//计算 m^a, O(loga), 本身没什么用, 注意这个按位处理的方法:-P
int exponent(int m,int a){
    int ret=1;
    for (;a;a>>=1,m*=m)
         if (a&1)
             ret*=m;
    return ret;
}
//计算幂取模 a^b mod n, O(logb)
int modular exponent(int a,int b,int n) { //a^b \mod n
    int ret=1;
    for (;b;b>>=1,a=(int)((i64)a)*a\%n)
         if (b&1)
             ret=(int)((i64)ret)*a%n;
    return ret;
}
//求解模线性方程 ax=b (mod n)
//返回解的个数,解保存在 sol[]中
```

```
//要求 n>0,解的范围 0..n-1
int modular_linear(int a,int b,int n,int* sol){
                   int d,e,x,y,i;
                   d=ext gcd(a,n,x,y);
                  if (b%d)
                                      return 0;
                   e=(x*(b/d)%n+n)%n;
                   for (i=0; i< d; i++)
                                      sol[i]=(e+i*(n/d))\%n;
                   return d;
}
//求解模线性方程组(中国余数定理)
// x = b[0] \pmod{w[0]}
// x = b[1] \pmod{w[1]}
// x = b[k-1] \pmod{w[k-1]}
//要求 w[i]>0,w[i]与 w[j]互质,解的范围 1..n,n=w[0]*w[1]*...*w[k-1]
int modular_linear_system(int b[],int w[],int k){
                   int d,x,y,a=0,m,n=1,i;
                   for (i=0; i< k; i++)
                                     n*=w[i];
                   for (i=0; i< k; i++)
                                     m=n/w[i];
                                      d=ext gcd(w[i],m,x,y);
                                      a=(a+y*m*b[i])%n;
                   }
                   return (a+n)%n;
}
4.3 素数
//用素数表判定素数,先调用 initprime
int plist[10000],pcount=0;
int prime(int n){
                   int i;
                   if ((n!=2\&\&!(n\%2))||(n!=3\&\&!(n\%3))||(n!=5\&\&!(n\%5))||(n!=7\&\&!(n\%7)))||(n!=3\&\&!(n\%3))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&!(n\%5))||(n!=3\&\&
                                      return 0;
                   for (i=0;plist[i]*plist[i] \le n;i++)
                                      if (!(n%plist[i]))
                                                        return 0;
                   return n>1;
```

}

```
void initprime(){
     int i;
     for (plist[pcount++]=2,i=3;i<50000;i++)
          if (prime(i))
              plist[pcount++]=i;
}
//miller rabin
//判断自然数 n 是否为素数
//time 越高失败概率越低,一般取 10 到 50
#include <stdlib.h>
#ifdef WIN32
typedef int64 i64;
#else
typedef long long i64;
#endif
int modular_exponent(int a,int b,int n){ //a^b mod n
     int ret;
     for (;b;b>>=1,a=(int)((i64)a)*a\%n)
          if (b&1)
              ret=(int)((i64)ret)*a\%n;
     return ret;
}
// Carmicheal number: 561,41041,825265,321197185
int miller rabin(int n,int time=10){
     if (n == 1 || (n! = 2 &\&! (n\%2)) || (n! = 3 &\&! (n\%3)) || (n! = 5 &\&! (n\%5)) || (n! = 7 &\&! (n\%7))) \\
          return 0;
     while (time--)
         if
(modular\_exponent(((rand()\&0x7fff<<16)+rand()\&0x7fff+rand()\&0x7fff)\%(n-1)+1,n-1,n)!=1)
              return 0;
     return 1;
}
4.4 欧拉函数
int gcd(int a,int b){
     return b?gcd(b,a%b):a;
}
inline int lcm(int a,int b){
```

### 5、数值计算

# 5.1 定积分计算(Romberg)

```
for (i=0; i<MAX N; i++) {
          R[0][i] = 0.0;
          R[1][i] = 0.0;
     }
    h = b-a;
     min = (int)(log(h*10.0)/log(2.0)); //h should be at most 0.1
     R[0][0] = ((*f)(a, l, t) + (*f)(b, l, t))*h*0.50;
    i = 1;
     temp2 = 1;
     while (i<MAX N){
          i++:
          R[1][0] = 0.0;
          for (j=1; j \le temp2; j++)
               R[1][0] += (*f)(a+h*((double)j-0.50), 1, t);
          R[1][0] = (R[0][0] + h*R[1][0])*0.50;
          temp4 = 4.0;
          for (j=1; j< i; j++) {
               R[1][j] = R[1][j-1] + (R[1][j-1]-R[0][j-1])/(temp4-1.0);
               temp4 *= 4.0;
          if ((fabs(R[1][i-1]-R[0][i-2])<eps)&&(i>min))
               return R[1][i-1];
          h *= 0.50;
          temp2 *= 2;
          for (j=0; j<i; j++)
               R[0][j] = R[1][j];
     }
     return R[1][MAX_N-1];
}
double Integral(double a, double b, double (*f)(double x, double y, double z), double eps,
                    double l, double t)
#define pi 3.1415926535897932
     int n;
     double R, p, res;
     n = (int)(floor)(b * t * 0.50 / pi);
     p = 2.0 * pi / t;
     res = b - (double)n * p;
    if (n)
          R = Romberg(a, p, f0, eps/(double)n, l, t);
     R = R * (double)n + Romberg(0.0, res, f0, eps, l, t);
```

```
return\ R/100.0; }
```

# 5.2 多项式求根(牛顿法)

```
/* 牛顿法解多项式的根
   输入: 多项式系数 c[], 多项式度数 n, 求在[a,b]间的根
   输出:根
   要求保证[a,b]间有根
double fabs( double x )
{
    return (x<0)? -x : x;
double f(int m, double c[], double x)
    int i;
    double p = c[m];
    for (i=m; i>0; i--)
         p = p*x + c[i-1];
    return p;
}
int newton(double x0, double *r,
             double c[], double cp[], int n,
             double a, double b, double eps)
{
    int MAX ITERATION = 1000;
    int i = 1;
    double x1, x2, fp, eps2 = eps/10.0;
    x1 = x0;
    while (i < MAX_ITERATION) {
         x2 = f(n, c, x1);
         fp = f(n-1, cp, x1);
         if ((fabs(fp)<0.000000001) && (fabs(x2)>1.0))
              return 0;
         x2 = x1 - x2/fp;
         if (fabs(x1-x2) \le eps2) {
              if (x2 \le a || x2 \ge b)
```

```
return 0;
               *r = x2;
               return 1;
          }
          x1 = x2;
          i++;
     return 0;
}
double Polynomial Root(double c[], int n, double a, double b, double eps)
{
     double *cp;
     int i;
     double root;
     cp = (double *)calloc(n, sizeof(double));
     for (i=n-1; i>=0; i--)
          cp[i] = (i+1)*c[i+1];
     }
     if (a>b) {
               root = a; a = b; b = root;
     if ((!newton(a, &root, c, cp, n, a, b, eps)) &&
          (!newton(b, &root, c, cp, n, a, b, eps)))
               newton((a+b)*0.5, &root, c, cp, n, a, b, eps);
     free(cp);
     if (fabs(root)<eps)
          return fabs(root);
     else
          return root;
}
```

# 5.3 周期性方程(追赶法)

```
/* 追赶法解周期性方程
   周期性方程定义: | a1 b1 c1 ...
                                                    = x1
                                                    = x2
                      a2 b2 c2 ...
                        . . .
                                                 X
                   cn−1 ...
                                                       xn-1
                                  an-1 bn-1
                  bn
                      cn
                                       an
                                                       xn
   输入: a[],b[],c[],x[]
   输出: 求解结果 X 在 x[]中
```

```
void run()
     c[0] /= b[0]; a[0] /= b[0]; x[0] /= b[0];
     for (int i = 1; i < N - 1; i +++) {
          double temp = b[i] - a[i] * c[i - 1];
          c[i] /= temp;
          x[i] = (x[i] - a[i] * x[i - 1]) / temp;
          a[i] = -a[i] * a[i - 1] / temp;
     }
     a[N-2] = -a[N-2] - c[N-2];
     for (int i = N - 3; i \ge 0; i - 1) {
          a[i] = -a[i] - c[i] * a[i + 1];
          x[i] = c[i] * x[i+1];
     }
     x[N-1] = (c[N-1] * x[0] + a[N-1] * x[N-2]);
     x[N-1] = (c[N-1] * a[0] + a[N-1] * a[N-2] + b[N-1]);
     for (int i = N - 2; i \ge 0; i - 1)
          x[i] += a[i] * x[N - 1];
}
```

### 6、图论─NP 搜索

#### 6.1 最大团

```
//最大团
//返回最大团大小和一个方案,传入图的大小 n 和邻接阵 mat
//mat[i][j]为布尔量
#define MAXN 60
void clique(int n, int* u, int mat[][MAXN], int size, int& max, int& bb, int* res, int* rr, int* c) {
    int i, j, vn, v[MAXN];
    if (n) {
         if (size + c[u[0]] \le max) return;
         for (i = 0; i < n + size - max && i < n; ++ i) {
              for (j = i + 1, vn = 0; j < n; ++ j)
                   if (mat[u[i]][u[j]])
                        v[vn ++] = u[j];
              rr[size] = u[i];
              clique(vn, v, mat, size + 1, max, bb, res, rr, c);
              if (bb) return;
         }
```

```
} else if (size > max) {
          max = size;
          for (i = 0; i < size; ++ i)
               res[i] = rr[i];
          bb = 1;
     }
}
int maxclique(int n, int mat[][MAXN], int *ret) {
     int max = 0, bb, c[MAXN], i, j;
     int vn, v[MAXN], rr[MAXN];
     for (c[i = n - 1] = 0; i >= 0; --i) {
          for (vn = 0, j = i + 1; j < n; ++ j)
               if (mat[i][j])
                    v[vn ++] = j;
          bb = 0:
          rr[0] = i;
          clique(vn, v, mat, 1, max, bb, ret, rr, c);
          c[i] = max;
     return max;
}
```

### 6.2 最大团(n<64)(faster)

```
/**

* WishingBone's ACM/ICPC Routine Library

*

* maximum clique solver

*/

#include <vector>

using std::vector;

// clique solver calculates both size and consitution of maximum clique

// uses bit operation to accelerate searching

// graph size limit is 63, the graph should be undirected

// can optimize to calculate on each component, and sort on vertex degrees

// can be used to solve maximum independent set

class clique {

public:

    static const long long ONE = 1;

    static const long long MASK = (1 << 21) - 1;
```

```
char* bits;
     int n, size, cmax[63];
     long long mask[63], cons;
     // initiate lookup table
     clique() {
          bits = new char[1 << 21];
          bits[0] = 0;
          for (int i = 1; i < 1 << 21; ++i) bits[i] = bits[i >> 1] + (i & 1);
     }
     ~clique() {
          delete bits:
    // search routine
     bool search(int step, int size, long long more, long long con);
    // solve maximum clique and return size
     int sizeClique(vector<vector<int>>& mat);
    // solve maximum clique and return constitution
     vector<int> consClique(vector<vector<int> >& mat);
};
// search routine
// step is node id, size is current solution, more is available mask, cons is
constitution mask
bool clique::search(int step, int size, long long more, long long cons) {
     if (step >= n) {
          // a new solution reached
          this->size = size;
          this->cons = cons;
          return true;
     long long now = ONE << step;
     if ((now \& more) > 0) {
          long long next = more & mask[step];
          if (size + bits[next & MASK] + bits[(next >> 21) & MASK] + bits[next >>
42] >= this->size
                    && size + cmax[step] \geq this-\geqsize) {
               // the current node is in the clique
               if (search(step + 1, size + 1, next, cons | now)) return true;
          }
     long long next = more & \simnow;
     if (size + bits[next & MASK] + bits[(next >> 21) & MASK] + bits[next >> 42]
> this->size) {
          // the current node is not in the clique
```

```
if (search(step + 1, size, next, cons)) return true;
     }
     return false;
}
// solve maximum clique and return size
int clique::sizeClique(vector<vector<int>>& mat) {
     n = mat.size();
     // generate mask vectors
     for (int i = 0; i < n; ++i) {
          mask[i] = 0;
          for (int j = 0; j < n; ++j) if (mat[i][j] > 0) mask[i] |= ONE << j;
     }
     size = 0;
     for (int i = n - 1; i \ge 0; --i) {
          search(i + 1, 1, mask[i], ONE << i);
          cmax[i] = size;
     }
     return size;
}
// solve maximum clique and return constitution
// calls sizeClique and restore cons
vector<int> clique::consClique(vector<vector<int> >& mat) {
     sizeClique(mat);
     vector<int> ret;
     for (int i = 0; i < n; ++i) if ((cons & (ONE << i)) > 0) ret.push_back(i);
     return ret;
}
```

#### 7、图论—连通性

### 7.1 无向图关键点(dfs 邻接阵)

```
//无向图的关键点,dfs 邻接阵形式,O(n^2)
//返回关键点个数,key[]返回点集
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
#define MAXN 110

void search(int n,int mat[][MAXN],int* dfn,int* low,int now,int& ret,int* key,int& cnt,int root,int& rd,int* bb){
    int i;
    dfn[now]=low[now]=++cnt;
```

```
for (i=0;i<n;i++)
          if (mat[now][i]){
                if (!dfn[i]) {
                     search(n,mat,dfn,low,i,ret,key,cnt,root,rd,bb);
                     if (low[i]<low[now])</pre>
                           low[now]=low[i];
                     if(low[i]>=dfn[now])
                           if (now!=root&&!bb[now])
                                 \text{key}[\text{ret}++]=\text{now},\text{bb}[\text{now}]=1;
                           else if(now==root)
                                rd++:
                }
                else if (dfn[i]<low[now])
                     low[now]=dfn[i];
          }
}
int key vertex(int n,int mat[][MAXN],int* key){
     int ret=0,i,cnt,rd,dfn[MAXN],low[MAXN],bb[MAXN];
     for (i=0;i< n;dfn[i++]=bb[i]=0);
     for (cnt=i=0;i< n;i++)
          if (!dfn[i]) {
                rd=0;
                search(n,mat,dfn,low,i,ret,key,cnt,i,rd,bb);
                if (rd>1&&!bb[i])
                     \text{key}[\text{ret++}]=i,\text{bb}[i]=1;
     return ret;
```

### 7.2 无向图关键边(dfs 邻接阵)

```
//无向图的关键边,dfs 邻接阵形式,O(n^2)
//返回关键边条数,key[][2]返回边集
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
#define MAXN 100

void search(int n,int mat[][MAXN],int* dfn,int* low,int now,int& cnt,int key[][2]){
    int i;
    for (low[now]=dfn[now],i=0;i<n;i++)
        if (mat[now][i]){
        if (!dfn[i]) {
            dfn[i]=dfn[now]+1;
```

```
search(n,mat,dfn,low,i,cnt,key);
                     if (low[i]>dfn[now])
                           \text{key}[\text{cnt}][0]=i,\text{key}[\text{cnt++}][1]=\text{now};
                     if (low[i]<low[now])</pre>
                           low[now]=low[i];
                else if (dfn[i]<dfn[now]-1&&dfn[i]<low[now])
                     low[now]=lev[i];
           }
}
int key edge(int n,int mat[][MAXN],int key[][2]){
     int ret=0,i,dfn[MAXN],low[MAXN];
     for (i=0;i< n;dfn[i++]=0);
     for (i=0;i<n;i++)
           if (!dfn[i])
                dfn[i]=1,bridge(n,mat,dfn,low,i,ret,key);
     return ret;
}
```

### 7.3 无向图的块(bfs 邻接阵)

```
//无向图的块,dfs 邻接阵形式,O(n^2)
//每产生一个块调用 dummy
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
#define MAXN 100
#include <iostream.h>
void dummy(int n,int* a){
    for (int i=0;i< n;i++)
         cout<<a[i]<<' ';
    cout << endl;
}
void search(int n,int mat[][MAXN],int* dfn,int* low,int now,int& cnt,int* st,int& sp){
    int i,m,a[MAXN];
    dfn[st[sp++]=now]=low[now]=++cnt;
    for (i=0; i< n; i++)
         if (mat[now][i]){
             if (!dfn[i]) {
                  search(n,mat,dfn,low,i,cnt,st,sp);
                  if (low[i] < low[now])</pre>
                       low[now]=low[i];
                  if(low[i] > = dfn[now]){
```

### 7.4 无向图连通分支(dfs/bfs 邻接阵)

```
//无向图连通分支,dfs 邻接阵形式,O(n^2)
//返回分支数,id 返回 1..分支数的值
//传入图的大小 n 和邻接阵 mat, 不相邻点边权 0
#define MAXN 100
void floodfill(int n,int mat[][MAXN],int* id,int now,int tag){
    int i;
    for (id[now]=tag,i=0;i< n;i++)
        if (!id[i]&&mat[now][i])
             floodfill(n,mat,id,i,tag);
}
int find components(int n,int mat[][MAXN],int* id){
    int ret,i;
    for (i=0;i< n;id[i++]=0);
    for (\text{ret}=i=0; i < n; i++)
        if (!id[i])
             floodfill(n,mat,id,i,++ret);
    return ret;
}
//无向图连通分支,bfs 邻接阵形式,O(n^2)
//返回分支数,id 返回 1..分支数的值
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
#define MAXN 100
```

```
\label{eq:components} \begin{tabular}{ll} int find\_components(int n,int mat[][MAXN],int* id) { & int ret,k,i,j,m; \\ for $(k=0;k< n;id[k++]=0);$ \\ for $(ret=k=0;k< n;k++)$ & if $(!id[k])$ & for $(id[k]=-1,ret++,m=1;m;)$ & for $(m=i=0;i< n;i++)$ & if $(id[i]==-1)$ & for $(m++,id[i]=ret,j=0;j< n;j++)$ & if $(!id[j]\&\&mat[i][j])$ & id[j]=-1; & return ret; \\ \end{tabular}
```

### 7.5 有向图强连通分支(dfs/bfs 邻接阵)

```
//有向图强连通分支,dfs 邻接阵形式,O(n^2)
//返回分支数,id 返回 1..分支数的值
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
#define MAXN 100
void search(int n,int mat[][MAXN],int* dfn,int* low,int now,int& cnt,int& tag,int* id,int* st,int&
sp){
    int i,j;
    dfn[st[sp++]=now]=low[now]=++cnt;
    for (i=0; i< n; i++)
         if (mat[now][i]){
              if (!dfn[i]) {
                  ssearch(n,mat,dfn,low,i,cnt,tag,id,st,sp);
                  if (low[i] < low[now])</pre>
                       low[now]=low[i];
              }
              else if (dfn[i]<dfn[now]){
                  for (j=0; j \le x \& st[j]!=i; j++);
                  if (j<cnt&&dfn[i]<low[now])
                       low[now]=dfn[i];
              }
    if (low[now]==dfn[now])
         for (tag++;st[sp]!=now;id[st[--sp]]=tag);
}
int find_components(int n,int mat[][MAXN],int* id){
    int ret=0,i,cnt,sp,st[MAXN],dfn[MAXN],low[MAXN];
```

```
for (i=0;i< n;dfn[i++]=0);
     for (sp=cnt=i=0;i< n;i++)
         if (!dfn[i])
              search(n,mat,dfn,low,i,cnt,ret,id,st,sp);
    return ret;
//有向图强连通分支,bfs 邻接阵形式,O(n^2)
//返回分支数,id 返回 1..分支数的值
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
#define MAXN 100
int find components(int n,int mat[][MAXN],int* id){
    int ret=0,a[MAXN],b[MAXN],c[MAXN],d[MAXN],i,j,k,t;
    for (k=0; k< n; id[k++]=0);
    for (k=0; k < n; k++)
         if (!id[k]){
              for (i=0;i<n;i++)
                  a[i]=b[i]=c[i]=d[i]=0;
              a[k]=b[k]=1;
              for (t=1;t;)
                  for (t=i=0;i< n;i++)
                       if (a[i]&&!c[i])
                            for (c[i]=t=1,j=0;j< n;j++)
                                 if (mat[i][j]&&!a[j])
                                     a[j]=1;
                       if (b[i]&&!d[i])
                            for (d[i]=t=1,j=0;j< n;j++)
                                 if (mat[j][i]&&!b[j])
                                     b[j]=1;
              for (ret++,i=0;i<n;i++)
                  if (a[i]&b[i])
                       id[i]=ret;
         }
    return ret;
}
```

#### 7.6 有向图最小点基(邻接阵)

```
//有向图最小点基,邻接阵形式,O(n^2)
//返回电集大小和点集
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
//需要调用强连通分支
#define MAXN 100
```

```
\label{eq:components} \begin{tabular}{ll} int base_vertex(int n,int mat[][MAXN],int* sets) & \\ int ret=0,id[MAXN],v[MAXN],i,j; \\ j=find_components(n,mat,id); \\ for (i=0;i<j;v[i++]=1); \\ for (i=0;i<n;i++) \\ for (j=0;j<n;j++) \\ if (id[i]!=id[j]&&mat[i][j]) \\ v[id[j]-1]=0; \\ for (i=0;i<n;i++) \\ if (v[id[i]-1]) \\ v[id[sets[ret++]=i]-1]=0; \\ return ret; \\ \end{tabular}
```

#### 8、图论—匹配

### 8.1 二分图最大匹配(hungary 邻接表)

```
//二分图最大匹配,hungary 算法,邻接表形式,复杂度 O(m*e)
//返回最大匹配数,传入二分图大小 m,n 和邻接表 list(只需一边)
//match1,match2 返回一个最大匹配,未匹配顶点 match 值为-1
#include <string.h>
#define MAXN 310
#define clr(x) memset(x,0xff,sizeof(int)*MAXN)
struct edge t{
    int from,to;
    edge_t* next;
};
int hungary(int m,int n,edge t* list[],int* match1,int* match2){
    int s[MAXN],t[MAXN],p,q,ret=0,i,j,k;edge t* e;
    for (clr(match1), clr(match2), i=0; i< m; ret+=(match1[i++]>=0))
         for ( clr(t), s[p=q=0]=i; p <= q \& match 1[i] < 0; p++)
             for (e=list[k=s[p]];e\&\&match1[i]<0;e=e->next)
                  if (t[j=e->to]<0){
                      s[++q]=match2[j],t[j]=k;
                      if (s[q] < 0)
                           for (p=j;p>=0;j=p)
                               match2[j]=k=t[j],p=match1[k],match1[k]=j;
                  }
    return ret;
}
```

### 8.2 二分图最大匹配(hungary 邻接阵)

```
//二分图最大匹配,hungary 算法,邻接阵形式,复杂度 O(m*m*n)
//返回最大匹配数、传入二分图大小 m,n 和邻接阵 mat, 非零元素表示有边
//match1,match2 返回一个最大匹配,未匹配顶点 match 值为-1
#include <string.h>
#define MAXN 310
#define clr(x) memset(x,0xff,sizeof(int)*MAXN)
int hungary(int m,int n,int mat[][MAXN],int* match1,int* match2){
    int s[MAXN],t[MAXN],p,q,ret=0,i,j,k;
    for (clr(match1), clr(match2), i=0; i< m; ret+=(match1[i++]>=0))
        for (clr(t),s[p=q=0]=i;p \le q \& match 1[i] \le 0;p++)
             for (k=s[p],j=0;j<n\&\&match1[i]<0;j++)
                 if (mat[k][j]\&\&t[j]<0){
                     s[++q]=match2[j],t[j]=k;
                     if (s[q]<0)
                          for (p=i;p>=0;j=p)
                              match2[j]=k=t[j],p=match1[k],match1[k]=j;
                 }
    return ret;
}
```

### 8.3 二分图最大匹配(hungary 正向表)

```
//二分图最大匹配,hungary 算法,正向表形式,复杂度 O(m*e)
//返回最大匹配数,传入二分图大小 m,n 和正向表 list,buf(只需一边)
//match1,match2 返回一个最大匹配,未匹配顶点 match 值为-1
#include <string.h>
#define MAXN 310
#define clr(x) memset(x,0xff,sizeof(int)*MAXN)
int hungary(int m,int n,int* list,int* buf,int* match1,int* match2){
    int s[MAXN],t[MAXN],p,q,ret=0,i,j,k,l;
    for (clr(match1), clr(match2), i=0; i< m; ret+=(match1[i++]>=0))
        for ( clr(t), s[p=q=0]=i; p <= q \& match 1[i] < 0; p++)
             for (l=list[k=s[p]]; l< list[k+1]&& match1[i]<0; l++)
                 if(t[j=buf[1]]<0){
                      s[++q]=match2[j],t[j]=k;
                      if (s[q]<0)
                          for (p=j;p>=0;j=p)
                              match2[j]=k=t[j],p=match1[k],match1[k]=j;
                 }
```

```
return ret;
```

#### 8.4 二分图最佳匹配(kuhn munkras 邻接阵)

```
//二分图最佳匹配,kuhn munkras 算法,邻接阵形式,复杂度 O(m*m*n)
//返回最佳匹配值,传入二分图大小 m,n 和邻接阵 mat,表示权值
//match1,match2 返回一个最佳匹配,未匹配顶点 match 值为-1
//一定注意 m<=n,否则循环无法终止
//最小权匹配可将权值取相反数
#include <string.h>
#define MAXN 310
#define inf 1000000000
#define clr(x) memset(x,0xff,sizeof(int)*n)
int kuhn munkras(int m,int n,int mat[][MAXN],int* match1,int* match2){
    int s[MAXN],t[MAXN],11[MAXN],12[MAXN],p,q,ret=0,i,j,k;
    for (i=0;i<m;i++)
         for (11[i]=-\inf,j=0;j< n;j++)
             11[i]=mat[i][j]>11[i]?mat[i][j]:11[i];
    for (i=0;i< n;12[i++]=0);
    for (clr(match1), clr(match2), i=0; i < m; i++) {
         for ( clr(t), s[p=q=0]=i; p <= q \& match 1[i] < 0; p++)
             for (k=s[p],j=0;j<n\&\&match1[i]<0;j++)
                  if (11[k]+12[j]==mat[k][j]\&\&t[j]<0){
                      s[++q]=match2[j],t[j]=k;
                      if (s[q]<0)
                           for (p=j;p>=0;j=p)
                               match2[j]=k=t[j],p=match1[k],match1[k]=j;
         if (match1[i]<0){
             for (i--,p=inf,k=0;k<=q;k++)
                  for (j=0; j< n; j++)
                      if (t[i]<0\&\&11[s[k]]+12[i]-mat[s[k]][i]<p)
                           p=11[s[k]]+12[j]-mat[s[k]][j];
             for (j=0;j< n;l2[j]+=t[j]<0?0:p,j++);
             for (k=0;k\leq q;11[s[k++]]=p);
         }
    }
    for (i=0; i < m; i++)
         ret+=mat[i][match1[i]];
    return ret;
```

## 8.5 一般图匹配(邻接表)

```
//一般图最大匹配,邻接表形式,复杂度 O(n*e)
//返回匹配顶点对数,match 返回匹配,未匹配顶点 match 值为-1
//传入图的顶点数 n 和邻接表 list
#define MAXN 100
struct edge t{
    int from,to;
    edge t* next;
};
int aug(int n,edge_t* list[],int* match,int* v,int now){
    int t,ret=0;edge t* e;
    v[now]=1;
    for (e=list[now];e;e=e->next)
         if(!v[t=e->to])
             if (match[t] < 0)
                  match[now]=t,match[t]=now,ret=1;
             else{
                  v[t]=1;
                  if (aug(n,list,match,v,match[t]))
                       match[now]=t,match[t]=now,ret=1;
                  v[t]=0;
             }
             if (ret)
                  break;
    v[now]=0;
    return ret;
}
int graph match(int n,edge t* list[],int* match){
    int v[MAXN],i,j;
    for (i=0;i<n;i++)
    v[i]=0,match[i]=-1;
    for (i=0,j=n;i<n\&\&j>=2;)
         if (match[i]<0&&aug(n,list,match,v,i))
             i=0, j=2;
         else
             i++;
    for (i=j=0;i< n;i++)
         j+=(match[i]>=0);
    return j/2;
}
```

## 8.6 一般图匹配(邻接阵)

```
//一般图最大匹配,邻接阵形式,复杂度 O(n^3)
//返回匹配顶点对数,match 返回匹配,未匹配顶点 match 值为-1
//传入图的顶点数 n 和邻接阵 mat
#define MAXN 100
int aug(int n,int mat[][MAXN],int* match,int* v,int now){
    int i,ret=0;
    v[now]=1;
    for (i=0;i<n;i++)
         if(!v[i]\&\&mat[now][i]){
             if (match[i]<0)
                  match[now]=i,match[i]=now,ret=1;
             else{
                 v[i]=1;
                 if (aug(n,mat,match,v,match[i]))
                      match[now]=i,match[i]=now,ret=1;
                 v[i]=0;
             }
             if (ret)
                 break;
         }
    v[now]=0;
    return ret;
}
int graph match(int n,int mat[][MAXN],int* match){
    int v[MAXN],i,j;
    for (i=0;i<n;i++)
    v[i]=0,match[i]=-1;
    for (i=0,j=n;i<n\&\&j>=2;)
         if (match[i]<0&&aug(n,mat,match,v,i))
             i=0, j=2;
         else
             i++;
    for (i=j=0;i< n;i++)
        j+=(match[i]>=0);
    return j/2;
}
```

### 8.7 一般图匹配(正向表)

```
//一般图最大匹配,正向表形式,复杂度 O(n*e)
//返回匹配顶点对数,match 返回匹配,未匹配顶点 match 值为-1
//传入图的顶点数 n 和正向表 list,buf
#define MAXN 100
int aug(int n,int* list,int* buf,int* match,int* v,int now){
    int i,t,ret=0;
    v[now]=1;
    for (i=list[now];i<list[now+1];i++)
         if (!v[t=buf[i]]){
              if (match[t] < 0)
                  match[now]=t,match[t]=now,ret=1;
              else{
                  v[t]=1;
                  if (aug(n,list,buf,match,v,match[t]))
                       match[now]=t,match[t]=now,ret=1;
                  v[t]=0;
              }
              if (ret)
                  break;
         }
    v[now]=0;
    return ret;
}
int graph match(int n,int* list,int* buf,int* match){
    int v[MAXN],i,j;
    for (i=0;i<n;i++)
    v[i]=0,match[i]=-1;
    for (i=0,j=n;i<n\&\&j>=2;)
         if (match[i]<0&&aug(n,list,buf,match,v,i))
              i=0, j=2;
         else
              i++;
    for (i=j=0;i< n;i++)
         j+=(match[i]>=0);
    return j/2;
}
```

#### 9、图论—网络流

### 9.1 最大流(邻接阵)

```
//返回最大流量.flow 返回每条边的流量
//传入网络节点数 n,容量 mat,源点 source,汇点 sink
#define MAXN 100
#define inf 1000000000
int max_flow(int n,int mat[][MAXN],int source,int sink,int flow[][MAXN]){
    int pre[MAXN],que[MAXN],d[MAXN],p,q,t,i,j;
    if (source==sink) return inf;
    for (i=0:i<n:i++)
        for (j=0; j< n; flow[i][j++]=0);
    for (;;){
        for (i=0;i< n;pre[i++]=0);
        pre[t=source]=source+1,d[t]=inf;
        for (p=q=0;p\leq q\&\&!pre[sink];t=que[p++])
             for (i=0;i<n;i++)
                 if(!pre[i]\&\&j=mat[t][i]-flow[t][i])
                     pre[que[q++]=i]=t+1,d[i]=d[t]<j?d[t]:j;
                 else if (!pre[i]&&j=flow[i][t])
                     pre[que[q++]=i]=-t-1,d[i]=d[t]<j?d[t]:j;
        if (!pre[sink]) break;
        for (i=sink;i!=source;)
             if (pre[i]>0)
                 flow[pre[i]-1][i]+=d[sink], i=pre[i]-1;
             else
                 flow[i][-pre[i]-1]-=d[sink],i=-pre[i]-1;
    for (j=i=0;i \le n;j+=flow[source][i++]);
    return j;
9.2 上下界最大流(邻接阵)
//求上下界网络最大流,邻接阵形式
//返回最大流量,-1 表示无可行流,flow 返回每条边的流量
//传入网络节点数 n,容量 mat,流量下界 bf,源点 source,汇点 sink
//MAXN 应比最大结点数多 2,无可行流返回-1 时 mat 未复原!
#define MAXN 100
#define inf 1000000000
int
     limit max flow(int
                               mat[][MAXN],int
                                                 bf[][MAXN],int
                                                                               sink,int
                         n,int
                                                                   source,int
flow[][MAXN]){
```

//求网络最大流,邻接阵形式

```
int i,j,sk,ks;
    if (source==sink) return inf;
     for (mat[n][n+1]=mat[n+1][n]=mat[n][n]=mat[n+1][n+1]=i=0;i< n;i++)
         for (mat[n][i]=mat[i][n]=mat[n+1][i]=mat[i][n+1]=i=0;j< n;j++)
              mat[i][j]=bf[i][j], mat[n][i]+=bf[j][i], mat[i][n+1]+=bf[i][j];
    sk=mat[source][sink],ks=mat[sink][source],mat[source][sink]=mat[sink][source]=inf;
     for (i=0;i< n+2;i++)
         for (j=0; j< n+2; flow[i][j++]=0);
     \max flow(n+2, mat, n, n+1, flow);
     for (i=0;i<n;i++)
         if (flow[n][i]<mat[n][i]) return -1;</pre>
     flow[source][sink]=flow[sink][source]=0,mat[source][sink]=sk,mat[sink][source]=ks;
     max_flow(n,mat,source,sink,flow);
    for (i=0; i< n; i++)
         for (j=0; j< n; j++)
              mat[i][j] += bf[i][j], flow[i][j] += bf[i][j];
     for (j=i=0;i< n;j+=flow[source][i++]);
    return i:
}
9.3 上下界最小流(邻接阵)
//求上下界网络最小流,邻接阵形式
//返回最大流量,-1 表示无可行流,flow 返回每条边的流量
//传入网络节点数 n,容量 mat,流量下界 bf,源点 source,汇点 sink
//MAXN 应比最大结点数多 2,无可行流返回-1 时 mat 未复原!
#define MAXN 100
#define inf 1000000000
     limit min flow(int
                           n,int
                                  mat[][MAXN],int
                                                       bf[][MAXN],int
                                                                         source, int
                                                                                      sink, int
flow[][MAXN]){
    int i,j,sk,ks;
    if (source==sink) return inf;
     for (mat[n][n+1]=mat[n+1][n]=mat[n][n]=mat[n+1][n+1]=i=0;i< n;i++)
         for (mat[n][i]=mat[i][n]=mat[n+1][i]=mat[i][n+1]=j=0;j< n;j++)
              mat[i][j] = bf[i][j], mat[n][i] + = bf[j][i], mat[i][n+1] + = bf[i][j];
    sk=mat[source][sink],ks=mat[sink][source],mat[source][sink]=mat[sink][source]=inf;
     for (i=0;i< n+2;i++)
         for (j=0; j< n+2; flow[i][j++]=0);
     \max flow(n+2, mat, n, n+1, flow);
     for (i=0;i< n;i++)
         if (flow[n][i]<mat[n][i]) return -1;</pre>
     flow[source][sink]=flow[sink][source]=0,mat[source][sink]=sk,mat[sink][source]=ks;
     _max_flow(n,mat,sink,source,flow);
```

```
\label{eq:for interpolation} \begin{split} & \text{for } (i=0; i < n; i++) \\ & \text{for } (j=0; j < n; j++) \\ & \text{mat}[i][j] += bf[i][j], flow[i][j] += bf[i][j]; \\ & \text{for } (j=i=0; i < n; j+=flow[source][i++]); \\ & \text{return } j; \\ \rbrace \end{split}
```

### 9.4 最大流无流量(邻接阵)

```
//求网络最大流,邻接阵形式
//返回最大流量
//传入网络节点数 n,容量 mat,源点 source,汇点 sink
//注意 mat 矩阵被修改
#define MAXN 100
#define inf 1000000000
int max flow(int n,int mat[][MAXN],int source,int sink){
    int v[MAXN],c[MAXN],p[MAXN],ret=0,i,j;
    for (;;){
         for (i=0;i<n;i++)
              v[i]=c[i]=0;
         for (c[source]=inf;;){
              for (j=-1,i=0;i< n;i++)
                  if(!v[i]\&\&c[i]\&\&(j=-1||c[i]>c[j]))
                       j=i;
              if (j<0) return ret;
              if (j==sink) break;
              for (v[j]=1,i=0;i< n;i++)
                  if (mat[j][i]>c[i]\&\&c[j]>c[i])
                       c[i]=mat[j][i]<c[j]?mat[j][i]:c[j],p[i]=j;
         for (ret+=j=c[i=sink];i!=source;i=p[i])
              mat[p[i]][i]=j,mat[i][p[i]]+=j;
}
```

### 9.5 最小费用最大流(邻接阵)

//求网络最小费用最大流,邻接阵形式 //返回最大流量,flow 返回每条边的流量,netcost 返回总费用 //传入网络节点数 n,容量 mat,单位费用 cost,源点 source,汇点 sink

```
int min cost max flow(int n,int mat[][MAXN],int cost[][MAXN],int source,int sink,int
flow[][MAXN],int& netcost){
     int pre[MAXN],min[MAXN],d[MAXN],i,j,t,tag;
     if (source==sink) return inf;
     for (i=0; i< n; i++)
          for (j=0; j< n; flow[i][j++]=0);
     for (netcost=0;;){
          for (i=0; i< n; i++)
               pre[i]=0,min[i]=inf;
          for (pre[source]=source+1,min[source]=0,d[source]=inf,tag=1;tag;)
               for (tag=t=0;t< n;t++)
                    if(d[t])
                         for (i=0; i< n; i++)
                              if (j=mat[t][i]-flow[t][i]\&\&min[t]+cost[t][i] < min[i])\\
                                   tag=1,min[i]=min[t]+cost[t][i],pre[i]=t+1,d[i]=d[t]<i?d[t]:j;
                              else if (j=flow[i][t]&&min[t]<inf&&min[t]-cost[i][t]<min[i])
                                   tag=1,min[i]=min[t]-cost[i][t],pre[i]=-t-1,d[i]=d[t]<j?d[t]:j;
          if (!pre[sink]) break;
          for (netcost+=min[sink]*d[i=sink];i!=source;)
               if (pre[i]>0)
                    flow[pre[i]-1][i]+=d[sink], i=pre[i]-1;
               else
                    flow[i][-pre[i]-1]=d[sink], i=-pre[i]-1;
     for (j=i=0;i < n;j+=flow[source][i++]);
     return j;
```

#### 10、 图论一应用

#### 10.1 欧拉回路(邻接阵)

```
//求欧拉回路或欧拉路,邻接阵形式,复杂度 O(n^2)
//返回路径长度,path 返回路径(有向图时得到的是反向路径)
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
//可以有自环与重边,分为无向图和有向图
#define MAXN 100
void find_path_u(int n,int mat[][MAXN],int now,int& step,int* path){
```

```
int i;
     for (i=n-1;i>=0;i--)
          while (mat[now][i]){
               mat[now][i]--,mat[i][now]--;
               find path u(n,mat,i,step,path);
     path[step++]=now;
}
void find path d(int n,int mat[][MAXN],int now,int& step,int* path){
     int i:
     for (i=n-1;i>=0;i--)
          while (mat[now][i]){
               mat[now][i]--;
               find_path_d(n,mat,i,step,path);
     path[step++]=now;
}
int euclid_path(int n,int mat[][MAXN],int start,int* path){
     int ret=0;
     find_path_u(n,mat,start,ret,path);
    find path d(n,mat,start,ret,path);
     return ret;
}
```

#### 10.2 树的前序表转化

```
}

void makepre(int n,node* list[],int* pre,int* map){
  int v[MAXN],id=0,i;
  for (i=0;i<n;v[i++]=0);
  prenode(n,list,pre,map,v,0,-1,id);
}</pre>
```

### 10.3 树的优化算法

```
//最大顶点独立集
int max_node_independent(int n,int* pre,int* set){
     int c[MAXN],i,ret=0;
     for (i=0;i<n;i++)
         c[i]=set[i]=0;
     for (i=n-1;i>=0;i--)
         if (!c[i]){
              set[i]=1;
              if (pre[i]!=-1)
                   c[pre[i]]=1;
              ret++;
     return ret;
}
//最大边独立集
int max_edge_independent(int n,int* pre,int* set){
     int c[MAXN],i,ret=0;
     for (i=0;i<n;i++)
         c[i]=set[i]=0;
     for (i=n-1;i>=0;i--)
         if (!c[i]&&pre[i]!=-1&&!c[pre[i]]){
              set[i]=1;
              c[pre[i]]=1;
              ret++;
     return ret;
}
//最小顶点覆盖集
int min_node_cover(int n,int* pre,int* set){
     int c[MAXN],i,ret=0;
     for (i=0;i<n;i++)
```

```
c[i]=set[i]=0;
     for (i=n-1;i>=0;i--)
          if (!c[i]&&pre[i]!=-1&&!c[pre[i]]){
               set[i]=1;
               c[pre[i]]=1;
               ret++;
     return ret;
}
//最小顶点支配集
int min node_dominant(int n,int* pre,int* set){
     int c[MAXN],i,ret=0;
     for (i=0; i< n; i++)
          c[i]=set[i]=0;
     for (i=n-1;i>=0;i--)
          if (!c[i]\&\&(pre[i]==-1||!set[pre[i]])){
               if(pre[i]!=-1){
                    set[pre[i]]=1;
                    c[pre[i]]=1;
                    if (pre[pre[i]]!=-1)
                         c[pre[pre[i]]]=1;
               }
               else
                    set[i]=1;
               ret++;
          }
     return ret;
}
```

### 10.4 拓扑排序(邻接阵)

```
//拓扑排序,邻接阵形式,复杂度 O(n^2)
//如果无法完成排序,返回 0,否则返回 1,ret 返回有序点列
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
#define MAXN 100

int toposort(int n,int mat[][MAXN],int* ret){
    int d[MAXN],i,j,k;
    for (i=0;i<n;i++)
        for (d[i]=j=0;j<n;d[i]+=mat[j++][i]);
    for (k=0;k<n;ret[k++]=i){
        for (i=0;d[i]&&i<n;i++);
        if (i==n)
```

```
return 0;
for (d[i]=-1,j=0;j<n;j++)
d[j]-=mat[i][j];
}
return 1;
}
```

#### 10.5 最佳边割集

```
//最佳边割集
#define MAXN 100
#define inf 1000000000
int max flow(int n,int mat[][MAXN],int source,int sink){
    int v[MAXN],c[MAXN],p[MAXN],ret=0,i,j;
    for (;;){
         for (i=0;i<n;i++)
              v[i]=c[i]=0;
         for (c[source]=inf;;){
               for (j=-1,i=0;i< n;i++)
                   if(!v[i]\&\&c[i]\&\&(j==-1||c[i]>c[j]))
              if (j<0) return ret;
              if (j==sink) break;
              for (v[j]=1,i=0;i < n;i++)
                   if(mat[j][i]>c[i]\&\&c[j]>c[i])
                        c[i]=mat[j][i]<c[j]?mat[j][i]:c[j],p[i]=j;
         for (ret+=j=c[i=sink];i!=source;i=p[i])
              mat[p[i]][i]=j, mat[i][p[i]]+=j;
}
int best edge cut(int n,int mat[][MAXN],int source,int sink,int set[][2],int& mincost){
    int m0[MAXN][MAXN],m[MAXN][MAXN],i,j,k,l,ret=0,last;
    if (source==sink)
         return -1;
    for (i=0; i< n; i++)
         for (j=0; j< n; j++)
              m0[i][j]=mat[i][j];
    for (i=0;i<n;i++)
         for (j=0;j<n;j++)
              m[i][j]=m0[i][j];
    mincost=last=max flow(n,m,source,sink);
```

#### 10.6 最佳点割集

```
//最佳顶点割集
#define MAXN 100
#define inf 1000000000
int max_flow(int n,int mat[][MAXN],int source,int sink){
     int v[MAXN],c[MAXN],p[MAXN],ret=0,i,j;
     for (;;){
          for (i=0;i<n;i++)
               v[i]=c[i]=0;
         for (c[source]=inf;;){
               for (j=-1,i=0;i< n;i++)
                    if(!v[i]\&\&c[i]\&\&(j==-1||c[i]>c[j]))
                        j=i;
               if (j<0) return ret;
               if (j==sink) break;
               for (v[j]=1,i=0;i< n;i++)
                   if(mat[j][i]>c[i]\&\&c[j]>c[i])
                         c[i]=mat[j][i]<c[j]?mat[j][i]:c[j],p[i]=j;
          for (ret+=j=c[i=sink];i!=source;i=p[i])
               mat[p[i]][i]=j,mat[i][p[i]]+=j;
}
```

int best\_vertex\_cut(int n,int mat[][MAXN],int\* cost,int source,int sink,int\* set,int& mincost){

```
int m0[MAXN][MAXN],m[MAXN][MAXN],i,j,k,ret=0,last;
if (source==sink||mat[source][sink])
    return -1;
for (i=0;i<n+n;i++)
    for (j=0; j< n+n; j++)
         m0[i][j]=0;
for (i=0; i< n; i++)
    for (j=0; j< n; j++)
         if (mat[i][j])
              m0[i][n+j]=inf;
for (i=0;i< n;i++)
    m0[n+i][i]=cost[i];
for (i=0;i<n+n;i++)
    for (j=0; j< n+n; j++)
         m[i][j]=m0[i][j];
mincost=last=max flow(n+n,m,source,n+sink);
for (k=0;k<n&&last;k++)
    if (k!=source&&k!=sink){
          for (i=0;i<n+n;i++)
              for (j=0;j< n+n;j++)
                   m[i][j]=m0[i][j];
         m[n+k][k]=0;
         if (max flow(n+n,m,source,n+sink)==last-cost[k]) {
              set[ret++]=k;
              m0[n+k][k]=0;
              last=cost[k];
         }
return ret;
```

#### 10.7 最小边割集

```
//最小边割集
#define MAXN 100
#define inf 1000000000

int max_flow(int n,int mat[][MAXN],int source,int sink){
    int v[MAXN],c[MAXN],p[MAXN],ret=0,i,j;
    for (;;){
        for (i=0;i<n;i++)
            v[i]=c[i]=0;
        for (c[source]=inf;;){
            for (j=-1,i=0;i<n;i++)
```

```
if(!v[i]\&\&c[i]\&\&(j=-1||c[i]>c[j]))
                         j=i;
               if (j<0) return ret;
               if (j==sink) break;
               for (v[j]=1,i=0;i< n;i++)
                    if (mat[j][i]>c[i]\&\&c[j]>c[i])
                         c[i]=mat[j][i]<c[j]?mat[j][i]:c[j],p[i]=j;
          }
          for (ret+=j=c[i=sink];i!=source;i=p[i])
               mat[p[i]][i]=j, mat[i][p[i]]+=j;
     }
}
int min_edge_cut(int n,int mat[][MAXN],int source,int sink,int set[][2]){
     int m0[MAXN][MAXN],m[MAXN][MAXN],i,j,k,l,ret=0,last;
     if (source==sink)
          return -1;
     for (i=0; i< n; i++)
          for (j=0; j< n; j++)
               m0[i][j]=(mat[i][j]!=0);
     for (i=0;i< n;i++)
          for (j=0; j< n; j++)
               m[i][j]=m0[i][j];
     last=max_flow(n,m,source,sink);
     for (k=0;k<n&&last;k++)
          for (l=0;l<n&&last;l++)
               if(m0[k][1]){
                    for (i=0;i<n+n;i++)
                         for (j=0;j< n+n;j++)
                              m[i][j]=m0[i][j];
                    m[k][1]=0;
                    if (max flow(n,m,source,sink)<last) {
                         set[ret][0]=k;
                         set[ret++][1]=l;
                         m0[k][1]=0;
                         last--;
                    }
     return ret;
}
```

### 10.8 最小点割集

//最小顶点割集

```
#define inf 1000000000
int max flow(int n,int mat[][MAXN],int source,int sink){
     int v[MAXN],c[MAXN],p[MAXN],ret=0,i,j;
     for (;;){
          for (i=0;i<n;i++)
               v[i]=c[i]=0;
          for (c[source]=inf;;){
               for (j=-1, i=0; i< n; i++)
                    if(!v[i]\&\&c[i]\&\&(j==-1||c[i]>c[j]))
                         j=i;
               if (j<0) return ret;
               if (j==sink) break;
               for (v[i]=1,i=0;i< n;i++)
                    if (mat[j][i]>c[i]\&\&c[j]>c[i])
                         c[i]=mat[j][i]<c[j]?mat[j][i]:c[j],p[i]=j;
          }
          for (ret+=j=c[i=sink];i!=source;i=p[i])
               mat[p[i]][i]=j, mat[i][p[i]]+=j;
     }
}
int min_vertex_cut(int n,int mat[][MAXN],int source,int sink,int* set){
     int m0[MAXN][MAXN],m[MAXN][MAXN],i,j,k,ret=0,last;
     if (source==sink||mat[source][sink])
          return -1;
     for (i=0;i<n+n;i++)
          for (j=0;j< n+n;j++)
               m0[i][j]=0;
     for (i=0;i<n;i++)
          for (j=0; j< n; j++)
               if(mat[i][j])
                    m0[i][n+j]=inf;
     for (i=0;i< n;i++)
          m0[n+i][i]=1;
     for (i=0;i<n+n;i++)
          for (j=0;j< n+n;j++)
               m[i][j]=m0[i][j];
     last=max flow(n+n,m,source,n+sink);
     for (k=0;k<n&&last;k++)
          if (k!=source&&k!=sink){
               for (i=0;i<n+n;i++)
                    for (j=0; j< n+n; j++)
```

#define MAXN 100

```
m[i][j]=m0[i][j];
m[n+k][k]=0;
if (max_flow(n+n,m,source,n+sink)<last){
    set[ret++]=k;
    m0[n+k][k]=0;
    last--;
}
return ret;
}</pre>
```

#### 10.9 最小路径覆盖

```
//最小路径覆盖,O(n^3)
//求解最小的路径覆盖图中所有点,有向图无向图均适用
//注意此问题等价二分图最大匹配,可以用邻接表或正向表减小复杂度
//返回最小路径条数,pre 返回前指针(起点-1),next 返回后指针(终点-1)
#include <string.h>
#define MAXN 310
#define clr(x) memset(x,0xff,sizeof(int)*n)
int hungary(int n,int mat[][MAXN],int* match1,int* match2){
    int s[MAXN],t[MAXN],p,q,ret=0,i,j,k;
    for (clr(match1), clr(match2),i=0;i<n;ret+=(match1[i++]>=0))
        for ( clr(t), s[p=q=0]=i; p <= q \& match 1[i] < 0; p++)
             for (k=s[p],j=0;j<n\&\&match1[i]<0;j++)
                 if (mat[k][j]\&\&t[j]<0){
                     s[++q]=match2[j],t[j]=k;
                     if (s[q]<0)
                         for (p=j;p>=0;j=p)
                             match2[j]=k=t[j],p=match1[k],match1[k]=j;
                 }
    return ret;
}
inline int path_cover(int n,int mat[][MAXN],int* pre,int* next){
    return n-hungary(n,mat,next,pre);
}
```

### 11、 图论—支撑树

### 11.1 最小生成树(kruskal 邻接表)

```
//无向图最小生成树,kruskal 算法,邻接表形式,复杂度 O(mlogm)
//返回最小生成树的长度、传入图的大小 n 和邻接表 list
//可更改边权的类型,edge[][2]返回树的构造,用边集表示
//如果图不连通.则对各连通分支构造最小生成树.返回总长度
#include <string.h>
#define MAXN 200
#define inf 1000000000
typedef double elem_t;
struct edge t{
    int from,to;
    elem t len;
    edge_t* next;
};
#define _ufind_run(x) for(;p[t=x];x=p[x],p[t]=(p[x]?p[x]:x))
#define run both ufind run(i); ufind run(j)
struct ufind{
    int p[MAXN],t;
    void init(){memset(p,0,sizeof(p));}
    void set_friend(int i,int j){_run_both;p[i]=(i==j?0:j);}
    int is friend(int i,int j){ run both;return i==j&&i;}
};
\#define _{cp(a,b)}((a).len<(b).len)
struct heap t{int a,b;elem t len;};
struct minheap{
    heap_t h[MAXN*MAXN];
    int n,p,c;
    void init(){n=0;}
    void ins(heap te){
        for (p=++n;p>1&& cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
        h[p]=e;
    }
    int del(heap_t& e){
        if (!n) return 0;
(e=h[p=1],c=2;c<n\&\&\_cp(h[c+=(c< n-1\&\&\_cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
        h[p]=h[n--];return 1;
};
elem_t kruskal(int n,edge_t* list[],int edge[][2]){
    ufind u;minheap h;
    edge_t* t;heap_t e;
```

### 11.2 最小生成树(kruskal 正向表)

```
//无向图最小生成树,kruskal 算法,正向表形式,复杂度 O(mlogm)
//返回最小生成树的长度,传入图的大小 n 和正向表 list,buf
//可更改边权的类型,edge[][2]返回树的构造,用边集表示
//如果图不连通,则对各连通分支构造最小生成树,返回总长度
#include <string.h>
#define MAXN 200
#define inf 1000000000
typedef double elem t;
struct edge t{
    int to;
    elem t len;
};
#define ufind_run(x) for(;p[t=x];x=p[x],p[t]=(p[x]?p[x]:x))
#define _run_both _ufind_run(i);_ufind_run(j)
struct ufind{
    int p[MAXN],t;
    void init(){memset(p,0,sizeof(p));}
    void set_friend(int i,int j){_run_both;p[i]=(i==j?0:j);}
    int is friend(int i,int j){ run both;return i==j&&i;}
};
#define cp(a,b) ((a).len<(b).len)
struct heap_t{int a,b;elem_t len;};
struct minheap{
    heap th[MAXN*MAXN];
    int n,p,c;
    void init(){n=0;}
    void ins(heap t e){
        for (p=++n;p>1&& cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
```

```
h[p]=e;
     }
     int del(heap t& e){
          if (!n) return 0;
          for
(e=h[p=1],c=2;c<n\&\& cp(h[c+=(c<n-1\&\& cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
          h[p]=h[n--]; return 1;
     }
};
elem t kruskal(int n,int* list,edge t* buf,int edge[][2]){
     ufind u;minheap h;
     heap_t e;elem_t ret=0;
     int i,j,m=0;
     u.init(),h.init();
     for (i=0;i< n;i++)
          for (j=list[i];j<list[i+1];j++)
               if (i<buf[j].to)
                    e.a=i,e.b=buf[j].to,e.len=buf[j].len,h.ins(e);
     while (m<n-1&&h.del(e))
          if (!u.is friend(e.a+1,e.b+1))
               edge[m][0]=e.a,edge[m][1]=e.b,ret+=e.len,u.set_friend(e.a+1,e.b+1);
     return ret;
}
```

## 11.3 最小生成树(prim+binary\_heap 邻接表)

```
//无向图最小生成树,prim 算法+二分堆,邻接表形式,复杂度 O(mlogm)
//返回最小生成树的长度,传入图的大小 n 和邻接表 list
//可更改边权的类型,pre[]返回树的构造,用父结点表示,根节点(第一个)pre 值为-1
//必须保证图的连通的!
#define MAXN 200
#define inf 1000000000
typedef double elem t;
struct edge t{
   int from,to;
   elem t len;
   edge_t* next;
};
#define cp(a,b)((a).d < (b).d)
struct heap_t{elem_t d;int v;};
struct heap{
   heap th[MAXN*MAXN];
```

```
int n,p,c;
     void init(){n=0;}
     void ins(heap te){
          for (p=++n;p>1&& cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
          h[p]=e;
     }
     int del(heap t& e){
          if (!n) return 0;
(e=h[p=1],c=2;c<n\&\& cp(h[c+=(c<n-1\&\& cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
          h[p]=h[n--]; return 1;
     }
};
elem t prim(int n,edge t* list[],int* pre){
     heap h;
     elem t min[MAXN],ret=0;
     edge t* t;heap t e;
     int v[MAXN],i;
     for (i=0;i<n;i++)
         min[i]=inf,v[i]=0,pre[i]=-1;
     h.init();e.v=0,e.d=0,h.ins(e);
     while (h.del(e))
          if (!v[e.v])
               for (v[e.v]=1,ret+=e.d,t=list[e.v];t;t=t->next)
                    if(!v[t->to]\&\&t->len<min[t->to])
                         pre[t->to]=t->from,min[e.v=t->to]=e.d=t->len,h.ins(e);
     return ret;
}
```

# 11.4 最小生成树(prim+binary\_heap 正向表)

```
//无向图最小生成树,prim 算法+二分堆,正向表形式,复杂度 O(mlogm)
//返回最小生成树的长度,传入图的大小 n 和正向表 list,buf
//可更改边权的类型,pre[]返回树的构造,用父结点表示,根节点(第一个)pre 值为-1
//必须保证图的连通的!
#define MAXN 200
#define inf 1000000000
typedef double elem_t;
struct edge_t{
    int to;
    elem_t len;
};
```

```
#define cp(a,b)((a).d < (b).d)
struct heap t{elem t d;int v;};
struct heap{
     heap th[MAXN*MAXN];
     int n,p,c;
     void init(){n=0;}
     void ins(heap te){
          for (p=++n;p>1&\&\_cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
          h[p]=e;
     }
     int del(heap t& e){
          if (!n) return 0;
(e=h[p=1],c=2;c<n&& cp(h[c+=(c<n-1&& cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
          h[p]=h[n--]; return 1;
};
elem t prim(int n,int* list,edge t* buf,int* pre){
     heap h;heap te;
     elem t min[MAXN],ret=0;
     int v[MAXN],i,j;
     for (i=0;i< n;i++)
          min[i]=inf,v[i]=0,pre[i]=-1;
     h.init();e.v=0,e.d=0,h.ins(e);
     while (h.del(e))
          if(!v[i=e.v])
               for (v[i]=1,ret+=e.d,j=list[i];j<list[i+1];j++)
                    if (!v[buf[j].to]&&buf[j].len<min[buf[j].to])</pre>
                         pre[buf[i].to]=i,min[e.v=buf[i].to]=e.d=buf[i].len,h.ins(e);
     return ret;
}
```

## 11.5 最小生成树(prim+mapped\_heap 邻接表)

```
//无向图最小生成树,prim 算法+映射二分堆,邻接表形式,复杂度 O(mlogn)
//返回最小生成树的长度,传入图的大小 n 和邻接表 list
//可更改边权的类型,pre[]返回树的构造,用父结点表示,根节点(第一个)pre 值为-1
//必须保证图的连通的!
#define MAXN 200
#define inf 1000000000
typedef double elem_t;
struct edge_t{
    int from,to;
```

```
elem_t len;
     edge_t* next;
};
#define cp(a,b)((a)<(b))
struct heap{
     elem_t h[MAXN+1];
     int ind[MAXN+1],map[MAXN+1],n,p,c;
     void init(){n=0;}
     void ins(int i,elem t e){
         for (p=++n;p>1\&\& cp(e,h[p>>1]);h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
         h[map[ind[p]=i]=p]=e;
     }
     int del(int i,elem_t& e){
         i=map[i]; if (i<1||i>n) return 0;
         for (e=h[p=i];p>1;h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
(c=2;c<n\&\&\_cp(h[c+=(c<n-1\&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<
=1);
         h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
     int delmin(int& i,elem_t& e){
         if (n<1) return 0;i=ind[1];
(e=h[p=1],c=2;c<n\&\& cp(h[c+=(c<n-1\&\& cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c]
],p=c,c<<=1);
         h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
     }
};
elem_t prim(int n,edge_t* list[],int* pre){
     heap h;
     elem_t min[MAXN],ret=0,e;
     edge_t* t;
     int v[MAXN],i;
     for (h.init(),i=0;i<n;i++)
         min[i]=(i?inf:0),v[i]=0,pre[i]=-1,h.ins(i,min[i]);
     while (h.delmin(i,e))
         for (v[i]=1,ret+=e,t=list[i];t;t=t->next)
              if(!v[t->to]\&\&t->len<min[t->to])
                   pre[t->to]=t->from,h.del(t->to,e),h.ins(t->to,min[t->to]=t->len);
     return ret;
}
```

### 11.6 最小生成树(prim+mapped\_heap 正向表)

```
//无向图最小生成树,prim 算法+映射二分堆,正向表形式,复杂度 O(mlogn)
//返回最小生成树的长度,传入图的大小 n 和正向表 list,buf
//可更改边权的类型,pre[]返回树的构造,用父结点表示,根节点(第一个)pre 值为-1
//必须保证图的连通的!
#define MAXN 200
#define inf 1000000000
typedef double elem t;
struct edge_t{
    int to;
    elem t len;
};
#define cp(a,b)((a)<(b))
struct heap{
    elem th[MAXN+1];
    int ind[MAXN+1],map[MAXN+1],n,p,c;
    void init(){n=0;}
    void ins(int i,elem t e){
        for (p=++n;p>1\&\& cp(e,h[p>>1]);h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
        h[map[ind[p]=i]=p]=e;
    }
    int del(int i,elem_t& e){
        i=map[i];if(i<1||i>n) return 0;
        for (e=h[p=i];p>1;h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
        for
(c=2;c<n\&\& cp(h[c+=(c<n-1\&\& cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<
=1);
        h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
    int delmin(int& i,elem t& e){
        if (n<1) return 0; i=ind[1];
(e=h[p=1],c=2;c<n\&\&\_cp(h[c+=(c< n-1\&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c]
],p=c,c<<=1);
        h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
    }
};
elem t prim(int n,int* list,edge t* buf,int* pre){
    elem t min[MAXN],ret=0,e;
    int v[MAXN],i,j;
```

# 11.7 最小生成树(prim 邻接阵)

```
//无向图最小生成树,prim 算法,邻接阵形式,复杂度 O(n^2)
//返回最小生成树的长度、传入图的大小 n 和邻接阵 mat,不相邻点边权 inf
//可更改边权的类型,pre[]返回树的构造,用父结点表示,根节点(第一个)pre 值为-1
//必须保证图的连通的!
#define MAXN 200
#define inf 1000000000
typedef double elem t;
elem t prim(int n,elem t mat[][MAXN],int* pre){
    elem t min[MAXN],ret=0;
    int v[MAXN],i,j,k;
    for (i=0;i<n;i++)
        min[i]=inf,v[i]=0,pre[i]=-1;
    for (\min[j=0]=0; j < n; j++){
        for (k=-1,i=0;i< n;i++)
            if(!v[i]&&(k==-1||min[i] < min[k]))
                k=i:
        for (v[k]=1,ret+=min[k],i=0;i< n;i++)
            if (!v[i]\&\&mat[k][i] < min[i])
                min[i]=mat[pre[i]=k][i];
    }
    return ret;
```

### 11.8 最小树形图(邻接阵)

```
//多源最小树形图,edmonds 算法,邻接阵形式,复杂度 O(n^3) //返回最小生成树的长度,构造失败返回负值 //传入图的大小 n 和邻接阵 mat,不相邻点边权 inf //可更改边权的类型,pre[]返回树的构造,用父结点表示 //传入时 pre[]数组清零,用-1 标出源点 #include <string.h>
```

```
#define MAXN 120
#define inf 1000000000
typedef int elem_t;
elem t edmonds(int n,elem t mat[][MAXN*2],int* pre){
     elem t ret=0;
     int c[MAXN*2][MAXN*2],l[MAXN*2],p[MAXN*2],m=n,t,i,j,k;
     for (i=0;i< n;l[i]=i,i++);
     do{
          memset(c,0,sizeof(c)),memset(p,0xff,sizeof(p));
          for (t=m,i=0;i<m;c[i][i]=1,i++);
          for (i=0;i<t;i++)
               if(l[i]==i\&\&pre[i]!=-1){
                    for (j=0; j < m; j++)
                         if(l[j]==j\&\&i!=j\&\&mat[j][i]<inf\&\&(p[i]==-1|mat[j][i]<mat[p[i]][i]))
                             p[i]=j;
                    if((pre[i]=p[i])==-1)
                         return -1;
                   if (c[i][p[i]]){
                         for (j=0;j \le m; mat[j][m] = mat[m][j] = inf,j++);
                         for (k=i; |[k]!=m; |[k]=m, k=p[k])
                              for (j=0;j< m;j++)
                                   if(l[j]==j){
                                        if (mat[j][k]-mat[p[k]][k] < mat[j][m])
                                             mat[j][m]=mat[j][k]-mat[p[k]][k];
                                        if (mat[k][j] < mat[m][j])
                                            mat[m][j]=mat[k][j];
                         c[m][m]=1,l[m]=m,m++;
                    for (j=0;j<m;j++)
                         if (c[i][j])
                              for (k=p[i];k!=-1\&\&l[k]==k;c[k][j]=1,k=p[k]);
               }
     while (t<m);
     for (;m-->n;pre[k]=pre[m])
          for (i=0;i<m;i++)
               if(||i||==m)
                    for (j=0;j<m;j++)
                         if(pre[j]==m\&\&mat[i][j]==mat[m][j])
                             pre[j]=i;
                    if (mat[pre[m]][m]==mat[pre[m]][i]-mat[pre[i]][i])
                         k=i;
```

```
}
for (i=0;i<n;i++)
    if (pre[i]!=-1)
        ret+=mat[pre[i]][i];
    return ret;
}</pre>
```

### 12、 图论─最短路径

# 12.1 最短路径(单源 bellman ford 邻接阵)

```
//单源最短路径,bellman ford 算法,邻接阵形式,复杂度 O(n^3)
//求出源 s 到所有点的最短路经、传入图的大小 n 和邻接阵 mat
//返回到各点最短距离 min[]和路径 pre[],pre[i]记录 s 到 i 路径上 i 的父结点,pre[s]=-1
//可更改路权类型,路权可为负,若图包含负环则求解失败,返回0
//优化:先删去负边使用 dijkstra 求出上界,加速迭代过程
#define MAXN 200
#define inf 1000000000
typedef int elem t;
int bellman ford(int n,elem t mat[][MAXN],int s,elem t* min,int* pre){
    int v[MAXN],i,j,k,tag;
    for (i=0; i< n; i++)
        min[i]=inf,v[i]=0,pre[i]=-1;
    for (\min[s]=0, j=0; j< n; j++){
        for (k=-1,i=0;i< n;i++)
            if (!v[i]\&\&(k==-1||min[i] < min[k]))
        for (v[k]=1,i=0;i < n;i++)
            if(!v[i]\&\&mat[k][i]>=0\&\&min[k]+mat[k][i]<min[i])
                 min[i]=min[k]+mat[pre[i]=k][i];
    for (tag=1,j=0;tag&&j<=n;j++)
        for (tag=i=0;i< n;i++)
            for (k=0; k< n; k++)
                if (min[k]+mat[k][i]<min[i])</pre>
                     min[i]=min[k]+mat[pre[i]=k][i],tag=1;
    return j<=n;
}
```

# 12.2 最短路径(单源 dijkstra+bfs 邻接表)

```
//单源最短路径,用于路权相等的情况,dijkstra 优化为 bfs,邻接表形式,复杂度 O(m)
//求出源 s 到所有点的最短路经、传入图的大小 n 和邻接表 list.边权值 len
//返回到各点最短距离 min[]和路径 pre[],pre[i]记录 s 到 i 路径上 i 的父结点,pre[s]=-1
//可更改路权类型,但必须非负且相等!
#define MAXN 200
#define inf 1000000000
typedef int elem t;
struct edge t{
    int from,to;
    edge t* next;
};
void dijkstra(int n,edge_t* list[],elem_t len,int s,elem_t* min,int* pre){
    edge t* t;
    int i,que[MAXN],f=0,r=0,p=1,l=1;
    for (i=0;i< n;i++)
        min[i]=inf;
    min[que[0]=s]=0,pre[s]=-1;
    for (;r \le f; l++, r=f+1, f=p-1)
        for (i=r;i<=f;i++)
            for (t=list[que[i]];t;t=t->next)
                if(min[t->to]==inf)
                     min[que[p++]=t->to]=len*l,pre[t->to]=que[i];
}
```

# 12.3 最短路径(单源 dijkstra+bfs 正向表)

```
//单源最短路径,用于路权相等的情况,dijkstra 优化为 bfs,正向表形式,复杂度 O(m)
//求出源 s 到所有点的最短路经、传入图的大小 n 和正向表 list,buf,边权值 len
//返回到各点最短距离 min[]和路径 pre[],pre[i]记录 s 到 i 路径上 i 的父结点,pre[s]=-1
//可更改路权类型,但必须非负且相等!
#define MAXN 200
#define inf 1000000000
typedef int elem t;
void dijkstra(int n,int* list,int* buf,elem t len,int s,elem t* min,int* pre){
    int i,que[MAXN],f=0,r=0,p=1,l=1,t;
    for (i=0; i< n; i++)
        min[i]=inf;
    min[que[0]=s]=0,pre[s]=-1;
    for (;r \le f;l++,r=f+1,f=p-1)
        for (i=r; i <= f; i++)
            for (t=list[que[i]];t<list[que[i]+1];t++)
                if (min[buf[t]] == inf)
```

```
min[que[p++]=buf[t]]=len*l,pre[buf[t]]=que[i];
```

# 12.4 最短路径(单源 dijkstra+binary\_heap 邻接表)

}

```
//单源最短路径,dijkstra 算法+二分堆,邻接表形式,复杂度 O(mlogm)
//求出源 s 到所有点的最短路经,传入图的大小 n 和邻接表 list
//返回到各点最短距离 min[]和路径 pre[],pre[i]记录 s 到 i 路径上 i 的父结点,pre[s]=-1
//可更改路权类型,但必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem t;
struct edge_t{
    int from,to;
    elem t len;
    edge t* next;
};
#define cp(a,b)((a).d < (b).d)
struct heap t{elem t d;int v;};
struct heap{
    heap_t h[MAXN*MAXN];
    int n,p,c;
    void init(){n=0;}
    void ins(heap t e){
        for (p=++n;p>1&\&\_cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
        h[p]=e;
    int del(heap t& e){
        if (!n) return 0;
(e=h[p=1],c=2;c<n\&\& cp(h[c+=(c<n-1\&\& cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
        h[p]=h[n--];return 1;
    }
};
void dijkstra(int n,edge t* list[],int s,elem t* min,int* pre){
    heap h;
    edge t* t;heap t e;
    int v[MAXN],i;
    for (i=0;i<n;i++)
        min[i]=inf,v[i]=0,pre[i]=-1;
    h.init();min[e.v=s]=e.d=0,h.ins(e);
    while (h.del(e))
```

```
 \begin{array}{c} if (!v[e.v]) \\ for (v[e.v]=1,t=list[e.v];t;t=t->next) \\ if (!v[t->to]&&min[t->from]+t->len<min[t->to]) \\ pre[t->to]=t->from,min[e.v=t->to]=e.d=min[t->from]+t->len,h.ins(e); \\ \end{array}
```

# 12.5 最短路径(单源 dijkstra+binary heap 正向表)

```
//单源最短路径,dijkstra 算法+二分堆,正向表形式,复杂度 O(mlogm)
//求出源 s 到所有点的最短路经,传入图的大小 n 和正向表 list,buf
//返回到各点最短距离 min[]和路径 pre[],pre[i]记录 s 到 i 路径上 i 的父结点,pre[s]=-1
//可更改路权类型,但必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem t;
struct edge t{
    int to;
    elem t len;
};
#define cp(a,b) ((a).d<(b).d)
struct heap_t{elem_t d;int v;};
struct heap{
    heap th[MAXN*MAXN];
    int n,p,c;
    void init(){n=0;}
    void ins(heap t e){
        for (p=++n;p>1&& cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
        h[p]=e;
    int del(heap t& e){
        if (!n) return 0;
        for
(e=h[p=1],c=2;c<n\&\& cp(h[c+=(c<n-1\&\& cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
        h[p]=h[n--]; return 1;
    }
};
void dijkstra(int n,int* list,edge t* buf,int s,elem t* min,int* pre){
    heap h;heap te;
    int v[MAXN],i,t,f;
    for (i=0; i< n; i++)
        min[i]=inf,v[i]=0,pre[i]=-1;
    h.init();min[e.v=s]=e.d=0,h.ins(e);
```

```
\label{eq:while} while (h.del(e)) $$ if (!v[e.v]) $$ for (v[f=e.v]=1,t=list[f];t<list[f+1];t++) $$ if (!v[buf[t].to]&&min[f]+buf[t].len<min[buf[t].to]) $$ pre[buf[t].to]=f,min[e.v=buf[t].to]=e.d=min[f]+buf[t].len,h.ins(e); $$ \}
```

# 12.6 最短路径(单源 dijkstra+mapped heap 邻接表)

```
//单源最短路径,dijkstra 算法+映射二分堆,邻接表形式,复杂度 O(mlogn)
//求出源 s 到所有点的最短路经,传入图的大小 n 和邻接表 list
//返回到各点最短距离 min[]和路径 pre[],pre[i]记录 s 到 i 路径上 i 的父结点,pre[s]=-1
//可更改路权类型.但必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem t;
struct edge_t{
    int from,to;
    elem t len;
    edge_t* next;
};
#define _cp(a,b) ((a)<(b))
struct heap{
    elem th[MAXN+1];
    int ind[MAXN+1],map[MAXN+1],n,p,c;
    void init()\{n=0;\}
    void ins(int i,elem t e){
        for (p=++n;p>1&& cp(e,h[p>>1]);h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
        h[map[ind[p]=i]=p]=e;
    }
    int del(int i,elem t& e){
        i=map[i];if(i<1||i>n) return 0;
        for (e=h[p=i];p>1;h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
(c=2;c<n\&\&\_cp(h[c+=(c<n-1\&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<
=1);
        h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
    int delmin(int& i,elem t& e){
        if (n<1) return 0; i=ind[1];
(e=h[p=1],c=2;c<n\&\&\_ep(h[c+=(c<n-1\&\&\_ep(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c]
],p=c,c<<=1);
```

# 12.7 最短路径(单源 dijkstra+mapped\_heap 正向表)

```
//单源最短路径,dijkstra 算法+映射二分堆,正向表形式,复杂度 O(mlogn)
//求出源 s 到所有点的最短路经,传入图的大小 n 和正向表 list,buf
//返回到各点最短距离 min[]和路径 pre[],pre[i]记录 s 到 i 路径上 i 的父结点,pre[s]=-1
//可更改路权类型.但必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem t;
struct edge_t{
    int to;
    elem t len;
};
\#define _{cp(a,b)}((a)<(b))
struct heap{
    elem_t h[MAXN+1];
    int ind[MAXN+1],map[MAXN+1],n,p,c;
    void init(){n=0;}
    void ins(int i,elem t e){
        for (p=++n;p>1\&\&\_cp(e,h[p>>1]);h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
        h[map[ind[p]=i]=p]=e;
    }
    int del(int i,elem t& e){
        i=map[i];if(i<1||i>n) return 0;
        for (e=h[p=i];p>1;h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
        for
(c=2;c<n\&\& cp(h[c+=(c<n-1\&\& cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<
```

```
=1);
          h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
     int delmin(int& i,elem t& e){
          if (n<1) return 0; i=ind[1];
(e=h[p=1],c=2;c<n\&\& cp(h[c+=(c<n-1\&\& cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c
],p=c,c<<=1);
          h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
     }
};
void dijkstra(int n,int* list,edge_t* buf,int s,elem_t* min,int* pre){
     heap h;elem t e;
     int v[MAXN],i,t;
     for (h.init(),i=0;i<n;i++)
          min[i]=((i==s)?0:inf),v[i]=0,pre[i]=-1,h.ins(i,min[i]);
     while (h.delmin(i,e))
          for (v[i]=1,t=list[i];t<list[i+1];t++)
               if (!v[buf[t].to]&&min[i]+buf[t].len<min[buf[t].to])</pre>
     pre[buf[t].to]=i,h.del(buf[t].to,e),min[buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to,e);
}
```

# 12.8 最短路径(单源 dijkstra 邻接阵)

```
//单源最短路径,dijkstra 算法,邻接阵形式,复杂度 O(n^2)
//求出源 s 到所有点的最短路经,传入图的顶点数 n,(有向)邻接矩阵 mat
//返回到各点最短距离 min[]和路径 pre[],pre[i]记录 s 到 i 路径上 i 的父结点,pre[s]=-1
//可更改路权类型,但必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem t;
void dijkstra(int n,elem t mat[][MAXN],int s,elem t* min,int* pre){
    int v[MAXN],i,j,k;
    for (i=0; i< n; i++)
        min[i]=inf,v[i]=0,pre[i]=-1;
    for (\min[s]=0, j=0, j< n, j++)
        for (k=-1,i=0;i< n;i++)
            if(!v[i]\&\&(k==-1||min[i]< min[k]))
                 k=i;
        for (v[k]=1,i=0;i< n;i++)
            if (!v[i]\&\&min[k]+mat[k][i]<min[i])
```

```
\label{eq:min[i]=min[k]+mat[pre[i]=k][i];} \\ \}
```

# 12.9 最短路径(多源 floyd\_warshall 邻接阵)

```
//多源最短路径,floyd warshall 算法,复杂度 O(n^3)
//求出所有点对之间的最短路经,传入图的大小和邻接阵
//返回各点间最短距离 min[]和路径 pre[],pre[i][j]记录 i 到 j 最短路径上 j 的父结点
//可更改路权类型,路权必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem t;
void floyd warshall(int n,elem t mat[][MAXN],elem t min[][MAXN],int pre[][MAXN]){
    int i,j,k;
    for (i=0; i< n; i++)
        for (j=0; j< n; j++)
            min[i][j]=mat[i][j],pre[i][j]=(i==j)?-1:i;
    for (k=0;k< n;k++)
        for (i=0;i<n;i++)
             for (j=0; j< n; j++)
                if (min[i][k]+min[k][j]<min[i][j])
                     min[i][j]=min[i][k]+min[k][j],pre[i][j]=pre[k][j];
}
```

# 13、 应用

# 13.1 Joseph 问题

```
// Joseph's Problem
// input: n,m -- the number of persons, the inteval between persons
// output: -- return the reference of last person
int josephus0(int n, int m)
{
    if (n == 2) return (m%2) ? 2 : 1;
    int v = (m+josephus0(n-1,m)) % n;
    if (v == 0) v = n;
    return v;
}
int josephus(int n, int m)
```

```
if (m == 1) return n;
if (n == 1) return 1;
if (m >= n) return josephus0(n,m);
int l = (n/m)*m;
int j = josephus(n - (n/m), m);
if (j <= n-l) return l+j;
j -= n-l;
int t = (j/(m-1))*m;
if ((j % (m-1)) == 0) return t-1;
return t + (j % (m-1));
}</pre>
```

### 13.2 N 皇后构造解

```
//N 皇后构造解,n>=4
void even1(int n,int *p){
     int i;
     for (i=1;i \le n/2;i++)
          p[i-1]=2*i;
     for (i=n/2+1; i \le n; i++)
          p[i-1]=2*i-n-1;
}
void even2(int n,int *p){
     int i;
     for (i=1;i \le n/2;i++)
          p[i-1]=(2*i+n/2-3)%n+1;
     for (i=n/2+1; i \le n; i++)
          p[i-1]=n-(2*(n-i+1)+n/2-3)%n;
}
void generate(int,int*);
void odd(int n,int *p){
     generate(n-1,p),p[n-1]=n;
}
void generate(int n,int *p){
     if (n&1)
          odd(n,p);
     else if (n\%6!=2)
          even1(n,p);
     else
```

```
even2(n,p); }
```

# 13.3 布尔母函数

```
//布尔母函数
//判 m[]个价值为 w[]的货币能否构成 value
//适合 m[]较大 w[]较小的情况
//返回布尔量
//传入货币种数 n,个数 m[],价值 w[]和目标值 value
#define MAXV 100000
int genfunc(int n,int* m,int* w,int value){
    int i,j,k,c;
    char r[MAXV];
    for (r[0]=i=1;i\le value;r[i++]=0);
    for (i=0;i< n;i++){
        for (j=0;j\leq w[i];j++)
             c=m[i]*r[k=j];
             while ((k+=w[i])<=value)
                 if(r[k])
                      c=m[i];
                 else if (c)
                     r[k]=1,c--;
             if (r[value])
                 return 1;
        }
    }
    return 0;
}
```

# 13.4 第 k 元素

```
//取第 k 个元素,k=0..n-1
//平均复杂度 O(n)
//注意 a[]中的顺序被改变

#define _cp(a,b) ((a)<(b))
typedef int elem_t;
elem_t kth_element(int n,elem_t* a,int k){
    elem_t t,key;
    int l=0,r=n-1,i,j;
    while (l<r){
```

### 13.5 幻方构造

```
//幻方构造(1!=2)
#define MAXN 100
void dllb(int l,int si,int sj,int sn,int d[][MAXN]){
     int n,i=0,j=1/2;
     for (n=1;n<=l*l;n++)
          d[i+si][j+sj]=n+sn;
          if (n%l){
               i=(i)?(i-1):(l-1);
               j=(j==l-1)?0:(j+1);
          }
          else
               i=(i==l-1)?0:(i+1);
     }
}
void magic odd(int l,int d[][MAXN]){
     dllb(1,0,0,0,d);
}
void magic_4k(int l,int d[][MAXN]){
     int i,j;
     for (i=0;i<1;i++)
          for (j=0;j<1;j++)
     d[i][j] = ((i\%4 = 0)|i\%4 = 3) & & (j\%4 = 0)|j\%4 = 3)|(i\%4 = 1)|i\%4 = 2) & & (j\%4 = 1)|j\%4 = 2)
?(l*l-(i*l+j)):(i*l+j+1);
}
void magic_other(int l,int d[][MAXN]){
     int i,j,t;
```

```
dllb(1/2,0,0,0,d);
     dllb(1/2,1/2,1/2,1*1/4,d);
     dllb(1/2,0,1/2,1*1/2,d);
     dllb(1/2,1/2,0,1*1/4*3,d);
     for (i=0;i<1/2;i++)
          for (j=0;j<1/4;j++)
                if (i!=1/4||i)
                     t=d[i][j],d[i][j]=d[i+1/2][j],d[i+1/2][j]=t;
     t=d[1/4][1/4],d[1/4][1/4]=d[1/4+1/2][1/4],d[1/4+1/2][1/4]=t;
     for (i=0;i<1/2;i++)
          for (j=l-1/4+1; j<1; j++)
                t=d[i][j],d[i][j]=d[i+1/2][j],d[i+1/2][j]=t;
}
void generate(int l,int d[][MAXN]){
     if (1%2)
          magic_odd(1,d);
     else if (1\%4 = = 0)
          magic 4k(1,d);
     else
          magic other(l,d);
}
```

# 13.6 模式匹配(kmp)

```
//模式匹配,kmp 算法,复杂度 O(m+n)
//返回匹配位置,-1表示匹配失败,传入匹配串和模式串和长度
//可更改元素类型,更换匹配函数
#define MAXN 10000
#define match(a,b) ((a)==(b))
typedef char elem t;
int pat_match(int ls,elem_t* str,int lp,elem_t* pat){
    int fail[MAXN]=\{-1\},i=0,j;
    for (j=1;j<1p;j++){
        for (i=fail[j-1];i\geq 0\&\&!_match(pat[i+1],pat[j]);i=fail[i]);
         fail[j]=(match(pat[i+1],pat[j])?i+1:-1);
    for (i=j=0;i<1s&&j<1p;i++)
         if (_match(str[i],pat[j]))
             j++;
         else if (j)
             j=fail[j-1]+1,i--;
    return j==lp?(i-lp):-1;
```

}

#### 13.7 逆序对数

```
//序列逆序对数,复杂度 O(nlogn)
//传入序列长度和内容,返回逆序对数
//可更改元素类型和比较函数
#include <string.h>
#define MAXN 1000000
#define cp(a,b)((a) \le (b))
typedef int elem t;
elem t tmp[MAXN];
int inv(int n,elem t* a){
    int l=n>>1,r=n-l,i,j;
    int ret=(r>1?(inv(1,a)+inv(r,a+1)):0);
    for (i=j=0;i<=1; tmp[i+j]=a[i],i++)
         for (ret+=j;j < r\&\&(i==l||! cp(a[i],a[l+j])); tmp[i+j]=a[l+j],j++);
    memcpy(a, tmp,sizeof(elem t)*n);
    return ret;
}
```

### 13.8 字符串最小表示

```
求字符串的最小表示
     输入:字符串
     返回:字符串最小表示的首字母位置(0...size-1)
template <class T>
int MinString(vector <T> &str)
{
     int i, j, k;
     vector \langle T \rangle ss(str.size() \langle < 1 \rangle;
     for (i = 0; i < str.size(); i ++) ss[i] = ss[i + str.size()] = str[i];
     for (i = k = 0, j = 1; k < str.size() && i < str.size() && j < str.size(); ) {
          for (k = 0; k < str.size() && ss[i + k] == ss[j + k]; k ++);
          if (k < str.size()) {
               if(ss[i+k] > ss[j+k])
                    i += k + 1;
               else j += k + 1;
               if (i == j) j ++;
```

```
\label{eq:continuous_problem} \left. \begin{array}{l} \text{return } i < j \ ? \ i : j; \\ \end{array} \right.
```

### 13.9 最长公共单调子序列

// 最长公共递增子序列, 时间复杂度 O(n^2 \* logn), 空间 O(n^2) \*n为a的大小,m为b的大小 \* 结果在 ans 中 \* "define\_cp(a,b) ((a)<(b))"求解最长严格递增序列 \*/ #define MAXN 1000 #define  $_{cp(a,b)}((a)<(b))$ typedef int elem\_t; elem\_t DP[MAXN][MAXN]; int num[MAXN], p[1<<20]; int LIS(int n, elem\_t \*a, int m, elem\_t \*b, elem\_t \*ans){ int i, j, l, r, k; DP[0][0] = 0;num[0] = (b[0] == a[0]);for(i = 1; i < m; i++) { num[i] = (b[i] == a[0]) || num[i-1];DP[i][0] = 0;} for(i = 1; i < n; i++){  $if(b[0] == a[i] &\& !num[0]) {$ num[0] = 1;DP[0][0] = i << 10;} for(j = 1; j < m; j++){ for(k=((l=0)+(r=num[j-1]-1))>>1; l<=r; k=(l+r)>>1)if(\_cp(a[DP[j-1][k]>>10], a[i])) 1=k+1; else r=k-1;  $if(1 < num[j-1] && i == (DP[j-1][1] >> 10)){$  $if(l \ge num[j]) DP[j][num[j]++] = DP[j-1][l];$ 

```
else DP[j][1] = _cp(a[DP[j][1] >> 10],a[i]) ? DP[j][1] : DP[j-1][1];
               }
               if(b[i] == a[i])
                    for(k=((l=0)+(r=num[j]-1))>>1; l<=r; k=(l+r)>>1)
                         if (cp(a[DP[j][k]>>10], a[i]))
                              1=k+1;
                         else
                              r=k-1;
                    DP[j][1] = (i << 10) + j;
                    num[j] += (l>=num[j]);
                    p[DP[j][1]] = 1 ? DP[j][1-1] : -1;
               }
         }
     }
     for (k=DP[m-1][i=num[m-1]-1];i>=0;ans[i--]=a[k>>10],k=p[k]);
     return num[m-1];
}
```

# 13.10 最长子序列

```
//最长单调子序列,复杂度 O(nlogn)
//注意最小序列覆盖和最长序列的对应关系,例如
//"define cp(a,b) ((a)>(b))"求解最长严格递减序列,则
//"define cp(a,b) (!((a)>(b)))"求解最小严格递减序列覆盖
//可更改元素类型和比较函数
#define MAXN 10000
#define cp(a,b)((a)>(b))
typedef int elem t;
int subseq(int n,elem t* a){
    int b[MAXN],i,l,r,m,ret=0;
    for (i=0;i< n;b[1]=i++,ret+=(1>ret))
        for (m=((l=1)+(r=ret))>>1; l<=r; m=(l+r)>>1)
            if (_cp(a[b[m]],a[i]))
                l=m+1;
            else
                r=m-1;
    return ret;
}
int subseq(int n,elem_t* a,elem_t* ans){
    int b[MAXN],p[MAXN],i,l,r,m,ret=0;
```

#### 13.11 最大子串匹配

```
//最大子串匹配,复杂度 O(mn)
//返回最大匹配值、传入两个串和串的长度、重载返回一个最大匹配
//注意做字符串匹配是串末的'\0'没有置!
//可更改元素类型,更换匹配函数和匹配价值函数
#include <string.h>
#define MAXN 100
#define \max(a,b) ((a)>(b)?(a):(b))
#define match(a,b) ((a)==(b))
#define value(a,b) 1
typedef char elem t;
int str_match(int m,elem_t* a,int n,elem_t* b){
    int match[MAXN+1][MAXN+1],i,j;
    memset(match,0,sizeof(match));
    for (i=0; i < m; i++)
         for (j=0; j< n; j++)
              match[i+1][j+1] = max(max(match[i][j+1], match[i+1][j]),
                                 (\text{match}[i][j] + \text{value}(a[i],b[i]))* \text{match}(a[i],b[j]);
    return match[m][n];
}
int str match(int m,elem t* a,int n,elem t* b,elem t* ret) {
    int match[MAXN+1][MAXN+1],last[MAXN+1][MAXN+1],i,j,t;
    memset(match,0,sizeof(match));
    for (i=0; i< m; i++)
         for (j=0; j< n; j++)
              match[i+1][j+1] = (match[i][j+1] > match[i+1][j]? match[i][j+1] : match[i+1][j]);
              last[i+1][j+1] = (match[i][j+1] > match[i+1][j]?3:1);
              if ((t=(match[i][j]+_value(a[i],b[i]))*_match(a[i],b[j]))>match[i+1][j+1])
                   match[i+1][j+1]=t, last[i+1][j+1]=2;
    for ( \text{:match}[i][j]; i=(\text{last}[t=i][j]>1), j=(\text{last}[t][j]<3) )
```

```
ret[match[i][j]-1] = (last[i][j] < 3?a[i-1]:b[j-1]); \\ return \ match[m][n]; \\ \}
```

### 13.12 最大子段和

```
//求最大子段和,复杂度 O(n)
//传入串长 n 和内容 list[]
//返回最大子段和,重载返回子段位置(maxsum=list[start]+...+list[end])
//可更改元素类型
typedef int elem t;
elem_t maxsum(int n,elem_t* list){
    elem t ret,sum=0;
    int i;
    for (ret=list[i=0]; i < n; i++)
         sum=(sum>0?sum:0)+list[i],ret=(sum>ret?sum:ret);
    return ret;
}
elem t maxsum(int n,elem t* list,int& start,int& end){
    elem_t ret,sum=0;
    int s,i;
    for (ret=list[start=end=s=i=0];i<n;i++,s=(sum>0?s:i))
         if ((sum=(sum>0?sum:0)+list[i])>ret)
             ret=sum,start=s,end=i;
    return ret;
}
```

# 13.13 最大子阵和

```
//求最大子阵和,复杂度 O(n^3)
//传入阵的大小 m,n 和内容 mat[][]
//返回最大子阵和,重载返回子阵位置(maxsum=list[s1][s2]+...+list[e1][e2])
//可更改元素类型
#define MAXN 100
typedef int elem_t;
elem_t maxsum(int m,int n,elem_t mat[][MAXN]){
    elem_t matsum[MAXN][MAXN+1],ret,sum;
    int i,j,k;
    for (i=0;i<m;i++)
        for (matsum[i][j=0]=0;j<n;j++)
```

```
matsum[i][j+1]=matsum[i][j]+mat[i][j];
    for (ret=mat[0][j=0]; j < n; j++)
         for (k=j;k\leq n;k++)
              for (sum=0,i=0;i<m;i++)
                   sum=(sum>0?sum:0)+matsum[i][k+1]-matsum[i][j],ret=(sum>ret?sum:ret);
    return ret;
}
elem t maxsum(int m,int n,elem t mat[][MAXN],int& s1,int& s2,int& e1,int& e2){
    elem t matsum[MAXN][MAXN+1],ret,sum;
    int i,j,k,s;
    for (i=0; i< m; i++)
         for (matsum[i][j=0]=0;j< n;j++)
              matsum[i][j+1]=matsum[i][j]+mat[i][j];
    for (ret=mat[s1=e1=0][s2=e2=j=0];j < n;j++)
         for (k=j;k\leq n;k++)
              for (sum=0,s=i=0;i< m;i++,s=(sum>0?s:i))
                   if((sum=(sum>0?sum:0)+matsum[i][k+1]-matsum[i][j])>ret)
                        ret=sum,s1=s,s2=i,e1=j,e2=k;
    return ret;
```

# 14、 其它

# 14.1 大数(只能处理正数)

```
#include <iostream.h>
#include <string.h>
#define DIGIT4
#define DEPTH
                   10000
#define MAX
                    100
typedef int bignum t[MAX+1];
int read(bignum t a,istream& is=cin){
    char buf[MAX*DIGIT+1],ch;
    memset((void*)a,0,sizeof(bignum_t));
    if (!(is>>buf)) return 0;
     for (a[0]=strlen(buf), i=a[0]/2-1; i>=0; i--)
         ch=buf[i],buf[i]=buf[a[0]-1-i],buf[a[0]-1-i]=ch;
    for (a[0]=(a[0]+DIGIT-1)/DIGIT, j=strlen(buf); j< a[0]*DIGIT; buf[j++]='0');
     for (i=1;i \le a[0];i++)
```

```
for (a[i]=0,j=0;j<DIGIT;j++)
               a[i]=a[i]*10+buf[i*DIGIT-1-j]-'0';
     for (;!a[a[0]]\&\&a[0]>1;a[0]--);
     return 1;
}
void write(const bignum t a,ostream& os=cout){
     int i,j;
     for (os << a[i=a[0]], i--; i; i--)
          for (j=DEPTH/10;j;j/=10)
               os < a[i]/j\%10;
}
int comp(const bignum t a,const bignum t b){
     int i;
     if (a[0]!=b[0])
          return a[0]-b[0];
     for (i=a[0];i;i--)
          if (a[i]!=b[i])
               return a[i]-b[i];
     return 0;
}
int comp(const bignum_t a,const int b){
     int c[12]=\{1\};
     for (c[1]=b;c[c[0]]>=DEPTH;c[c[0]+1]=c[c[0]]/DEPTH,c[c[0]]%=DEPTH,c[0]++);
     return comp(a,c);
}
int comp(const bignum t a,const int c,const int d,const bignum t b){
     int i,t=0,O=-DEPTH*2;
     if (b[0]-a[0] < d\&\&c)
          return 1;
     for (i=b[0];i>d;i--){
          t=t*DEPTH+a[i-d]*c-b[i];
          if (t>0) return 1;
          if (t<O) return 0;
     }
     for (i=d;i;i--)
          t=t*DEPTH-b[i];
          if (t>0) return 1;
          if (t<O) return 0;
     return t>0;
```

```
}
void add(bignum_t a,const bignum_t b){
    int i;
    for (i=1;i \le b[0];i++)
         if ((a[i]+=b[i])>=DEPTH)
              a[i]-=DEPTH,a[i+1]++;
    if (b[0] > = a[0])
         a[0]=b[0];
    else
         for (;a[i] \ge DEPTH\&\&i < a[0];a[i] = DEPTH,i++,a[i]++);
    a[0]+=(a[a[0]+1]>0);
}
void add(bignum_t a,const int b){
    int i=1:
    for (a[1]+=b;a[i]>=DEPTH\&\&i< a[0];a[i+1]+=a[i]/DEPTH,a[i]%=DEPTH,i++);
    for (a[a[0]] = DEPTH; a[a[0]+1] = a[a[0]]/DEPTH, a[a[0]]\% = DEPTH, a[0]++);
}
void sub(bignum t a,const bignum t b){
    int i;
     for (i=1;i \le b[0];i++)
         if ((a[i]-=b[i])<0)
              a[i+1]--,a[i]+=DEPTH;
     for (;a[i]<0;a[i]+=DEPTH,i++,a[i]--);
    for (;!a[a[0]]\&\&a[0]>1;a[0]--);
}
void sub(bignum t a,const int b){
    int i=1;
    for
(a[1]-b;a[i]<0;a[i+1]+=(a[i]-DEPTH+1)/DEPTH,a[i]-=(a[i]-DEPTH+1)/DEPTH*DEPTH,i++);
    for (;!a[a[0]]\&\&a[0]>1;a[0]--);
}
void sub(bignum_t a,const bignum_t b,const int c,const int d){
    int i,O=b[0]+d;
    for (i=1+d;i<=O;i++)
         if((a[i]-=b[i-d]*c)<0)
              a[i+1]+=(a[i]-DEPTH+1)/DEPTH,a[i]-=(a[i]-DEPTH+1)/DEPTH*DEPTH;
     for (;a[i]<0;a[i+1]+=(a[i]-DEPTH+1)/DEPTH,a[i]=(a[i]-DEPTH+1)/DEPTH*DEPTH,i++);
    for (;!a[a[0]]\&\&a[0]>1;a[0]--);
```

```
void mul(bignum t c,const bignum t a,const bignum t b){
    int i,j;
    memset((void*)c,0,sizeof(bignum t));
    for (c[0]=a[0]+b[0]-1,i=1;i \le a[0];i++)
         for (j=1;j<=b[0];j++)
              if((c[i+j-1]+=a[i]*b[j])>=DEPTH)
                   c[i+j]+=c[i+j-1]/DEPTH, c[i+j-1]%=DEPTH;
    for (c[0]+=(c[c[0]+1]>0); !c[c[0]] \&\&c[0]>1; c[0]--);
}
void mul(bignum t a,const int b){
    int i;
    for (a[1]*=b,i=2;i\leq=a[0];i++)
         a[i]*=b;
         if (a[i-1] > = DEPTH)
              a[i]+=a[i-1]/DEPTH, a[i-1]\%=DEPTH;
    }
     for (a[a[0]] \ge DEPTH, a[a[0]+1] = a[a[0]]/DEPTH, a[a[0]] = DEPTH, a[0]++);
    for (;!a[a[0]]\&\&a[0]>1;a[0]--);
}
void mul(bignum t b,const bignum t a,const int c,const int d){
    int i;
    memset((void*)b,0,sizeof(bignum t));
    for (b[0]=a[0]+d,i=d+1;i\leq=b[0];i++)
         if((b[i]+=a[i-d]*c)>=DEPTH)
              b[i+1]+=b[i]/DEPTH, b[i]\%=DEPTH;
     for (;b[b[0]+1];b[0]++,b[b[0]+1]=b[b[0]]/DEPTH,b[b[0]]\%=DEPTH);
    for (;!b[b[0]]\&\&b[0]>1;b[0]--);
}
void div(bignum_t c,bignum_t a,const bignum_t b){
    int h,l,m,i;
    memset((void*)c,0,sizeof(bignum t));
    c[0]=(b[0]< a[0]+1)?(a[0]-b[0]+2):1;
     for (i=c[0];i;sub(a,b,c[i]=m,i-1),i--)
         for (h=DEPTH-1,l=0,m=(h+l+1)>>1;h>l;m=(h+l+1)>>1)
              if (comp(b,m,i-1,a)) h=m-1;
              else l=m;
    for (; |c[c[0]] \& \&c[0] > 1; c[0] --);
    c[0]=c[0]>1?c[0]:1;
}
```

```
void div(bignum t a,const int b,int& c){
     int i:
     for (c=0,i=a[0];i;c=c*DEPTH+a[i],a[i]=c/b,c%=b,i--);
     for (;!a[a[0]]\&\&a[0]>1;a[0]--);
}
void sqrt(bignum t b,bignum t a){
     int h,l,m,i;
     memset((void*)b,0,sizeof(bignum t));
     for (i=b[0]=(a[0]+1)>>1;i;sub(a,b,m,i-1),b[i]+=m,i--)
          for (h=DEPTH-1,l=0,b[i]=m=(h+l+1)>>1;h>l;b[i]=m=(h+l+1)>>1)
               if (comp(b,m,i-1,a)) h=m-1;
               else l=m;
     for (;!b[b[0]]\&\&b[0]>1;b[0]--);
     for (i=1;i \le b[0];b[i++] >>=1);
}
int length(const bignum t a){
     int t,ret;
     for (ret=(a[0]-1)*DIGIT,t=a[a[0]];t;t/=10,ret++);
     return ret>0?ret:1;
}
int digit(const bignum_t a,const int b){
     int i,ret;
     for (ret=a[(b-1)/DIGIT+1],i=(b-1)%DIGIT;i;ret/=10,i--);
     return ret%10;
}
int zeronum(const bignum ta){
     int ret,t;
     for (ret=0;!a[ret+1];ret++);
     for (t=a[ret+1],ret*=DIGIT;!(t%10);t/=10,ret++);
     return ret;
}
void comp(int* a,const int l,const int h,const int d){
     int i,j,t;
     for (i=1;i<=h;i++)
          for (t=i,j=2;t>1;j++)
               while (!(t%j))
                    a[j]+=d,t/=j;
}
```

```
void convert(int* a,const int h,bignum t b){
     int i,j,t=1;
     memset(b,0,sizeof(bignum t));
     for (b[0]=b[1]=1,i=2;i<=h;i++)
          if (a[i])
               for (j=a[i];j;t*=i,j--)
                    if (t*i>DEPTH)
                         mul(b,t),t=1;
     mul(b,t);
}
void combination(bignum t a,int m,int n){
     int* t=new int[m+1];
     memset((void*)t,0,sizeof(int)*(m+1));
     comp(t,n+1,m,1);
     comp(t,2,m-n,-1);
     convert(t,m,a);
    delete ∏t;
}
void permutation(bignum t a,int m,int n){
     int i,t=1;
     memset(a,0,sizeof(bignum t));
     a[0]=a[1]=1;
     for (i=m-n+1; i \le m; t*=i++)
          if (t*i>DEPTH)
               mul(a,t),t=1;
     mul(a,t);
}
#define SGN(x) ((x)>0?1:((x)<0?-1:0))
#define ABS(x) ((x)>0?(x):-(x))
int read(bignum_t a,int &sgn,istream& is=cin){
     char str[MAX*DIGIT+2],ch,*buf;
     int i,j;
     memset((void*)a,0,sizeof(bignum_t));
     if (!(is>>str)) return 0;
     buf=str,sgn=1;
     if (*buf=='-') sgn=-1,buf++;
     for (a[0]=strlen(buf), i=a[0]/2-1; i>=0; i--)
          ch=buf[i],buf[i]=buf[a[0]-1-i],buf[a[0]-1-i]=ch;
     for (a[0]=(a[0]+DIGIT-1)/DIGIT, j=strlen(buf); j< a[0]*DIGIT; buf[j++]='0');
     for (i=1;i \le a[0];i++)
```

```
for \ (a[i]=0,j=0,j<DIGIT,j++) \\ a[i]=a[i]*10+buf[i*DIGIT-1-j]-'0'; \\ for \ (;!a[a[0]]\&\&a[0]>1;a[0]--); \\ if \ (a[0]==1\&\&!a[1]) \ sgn=0; \\ return \ 1; \\ \}
```

# 14.2 分数

```
struct frac{
     int num, den;
};
double fabs(double x){
     return x>0?x:-x;
int gcd(int a,int b){
     int t;
     if (a<0)
          a=-a;
     if (b<0)
          b=-b;
    if (!b)
          return a;
     while (t=a%b)
          a=b,b=t;
     return b;
}
void simplify(frac& f){
     int t;
     if (t=gcd(f.num,f.den))
          f.num/=t,f.den/=t;
     else
          f.den=1;
}
frac f(int n,int d,int s=1){
     frac ret;
     if (d<0)
          ret.num=-n,ret.den=-d;
     else
          ret.num=n,ret.den=d;
```

```
if (s)
          simplify(ret);
     return ret;
}
frac convert(double x){
     frac ret;
     for (ret.den=1; fabs(x-int(x)) > 1e-10; ret.den*=10,x*=10);
     ret.num=(int)x;
     simplify(ret);
     return ret;
}
int fragemp(frac a,frac b){
     int g1=gcd(a.den,b.den),g2=gcd(a.num,b.num);
     if (!g1||!g2)
          return 0;
    return b.den/g1*(a.num/g2)-a.den/g1*(b.num/g2);
}
frac add(frac a,frac b){
     int g1=gcd(a.den,b.den),g2,t;
    if (!g1)
          return f(1,0,0);
     t=b.den/g1*a.num+a.den/g1*b.num;
     g2=gcd(g1,t);
     return f(t/g2,a.den/g1*(b.den/g2),0);
}
frac sub(frac a,frac b){
     return add(a,f(-b.num,b.den,0));
}
frac mul(frac a,frac b){
     int t1=gcd(a.den,b.num),t2=gcd(a.num,b.den);
    if (!t1||!t2)
          return f(1,1,0);
    return f(a.num/t2*(b.num/t1),a.den/t1*(b.den/t2),0);
}
frac div(frac a,frac b){
     return mul(a,f(b.den,b.num,0));
}
```

#### 14.3 矩阵

```
define MAXN 100
#define fabs(x) ((x)>0?(x):-(x))
#define zero(x) (fabs(x)<1e-10)
struct mat{
     int n,m;
     double data[MAXN][MAXN];
};
int mul(mat& c,const mat& a,const mat& b){
     int i,j,k;
     if (a.m!=b.n)
          return 0;
     c.n=a.n,c.m=b.m;
     for (i=0; i< c.n; i++)
          for (j=0;j< c.m;j++)
               for (c.data[i][j]=k=0;k< a.m;k++)
                    c.data[i][j]+=a.data[i][k]*b.data[k][j];
     return 1;
}
int inv(mat& a){
     int i,j,k,is[MAXN],js[MAXN];
     double t;
     if (a.n!=a.m)
          return 0;
     for (k=0; k< a.n; k++){
          for (t=0,i=k;i<a.n;i++)
               for (j=k;j<a.n;j++)
                    if (fabs(a.data[i][j])>t)
                          t=fabs(a.data[is[k]=i][js[k]=j]);
          if (zero(t))
               return 0;
          if (is[k]!=k)
               for (j=0; j< a.n; j++)
                    t=a.data[k][j],a.data[k][j]=a.data[is[k]][j],a.data[is[k]][j]=t;
          if(js[k]!=k)
               for (i=0;i<a.n;i++)
                    t=a.data[i][k],a.data[i][k]=a.data[i][js[k]],a.data[i][js[k]]=t;
          a.data[k][k]=1/a.data[k][k];
          for (j=0; j< a.n; j++)
```

```
if(j!=k)
                     a.data[k][j]*=a.data[k][k];
          for (i=0; i< a.n; i++)
               if(i!=k)
                     for (j=0; j< a.n; j++)
                          if(j!=k)
                               a.data[i][j]=a.data[i][k]*a.data[k][j];
          for (i=0;i<a.n;i++)
               if(i!=k)
                     a.data[i][k]*=-a.data[k][k];
     }
     for (k=a.n-1;k>=0;k--)
          for (j=0; j< a.n; j++)
               if (js[k]!=k)
                     t=a.data[k][j],a.data[k][j]=a.data[js[k]][j],a.data[js[k]][j]=t;
          for (i=0;i<a.n;i++)
               if(is[k]!=k)
                     t=a.data[i][k],a.data[i][k]=a.data[i][is[k]],a.data[i][is[k]]=t;
     }
     return 1;
}
double det(const mat& a){
     int i,j,k,sign=0;
     double b[MAXN][MAXN],ret=1,t;
     if (a.n!=a.m)
          return 0;
     for (i=0; i< a.n; i++)
          for (j=0;j<a.m;j++)
               b[i][j]=a.data[i][j];
     for (i=0; i< a.n; i++)
          if (zero(b[i][i])){
                for (j=i+1; j< a.n; j++)
                     if (!zero(b[j][i]))
                          break;
               if(j==a.n)
                     return 0;
               for (k=i;k<a.n;k++)
                     t=b[i][k],b[i][k]=b[j][k],b[j][k]=t;
               sign++;
          }
          ret*=b[i][i];
          for (k=i+1;k< a.n;k++)
               b[i][k]/=b[i][i];
```

# 14.4 线性方程组

```
#define MAXN 100
#define fabs(x) ((x)>0?(x):-(x))
#define eps 1e-10
//列主元 gauss 消去求解 a[][]x[]=b[]
//返回是否有唯一解,若有解在 b[]中
int gauss cpivot(int n,double a[][MAXN],double b[]){
     int i,j,k,row;
     double maxp,t;
     for (k=0;k< n;k++){
          for (\max p=0, i=k; i < n; i++)
              if (fabs(a[i][k])>fabs(maxp))
                   maxp=a[row=i][k];
          if (fabs(maxp)<eps)
              return 0;
          if(row!=k){
              for (j=k; j < n; j++)
                   t=a[k][j],a[k][j]=a[row][j],a[row][j]=t;
              t=b[k],b[k]=b[row],b[row]=t;
          for (j=k+1; j< n; j++)
              a[k][j]/=maxp;
              for (i=k+1;i< n;i++)
                   a[i][j] = a[i][k] * a[k][j];
          b[k]/=maxp;
          for (i=k+1;i<n;i++)
              b[i]=b[k]*a[i][k];
     for (i=n-1;i>=0;i--)
          for (j=i+1; j< n; j++)
              b[i]=a[i][j]*b[j];
     return 1;
```

```
}
//全主元 gauss 消去解 a[][]x[]=b[]
//返回是否有唯一解,若有解在 b[]中
int gauss tpivot(int n,double a[][MAXN],double b[]){
     int i,j,k,row,col,index[MAXN];
     double maxp,t;
     for (i=0;i<n;i++)
          index[i]=i;
     for (k=0;k< n;k++)
          for (maxp=0,i=k;i<n;i++)
               for (j=k;j\leq n;j++)
                   if (fabs(a[i][j])>fabs(maxp))
                         maxp=a[row=i][col=j];
          if (fabs(maxp)<eps)
               return 0;
          if(col!=k){
               for (i=0; i< n; i++)
                   t=a[i][col],a[i][col]=a[i][k],a[i][k]=t;
              j=index[col],index[col]=index[k],index[k]=j;
          if(row!=k){
               for (j=k; j < n; j++)
                    t=a[k][j],a[k][j]=a[row][j],a[row][j]=t;
               t=b[k],b[k]=b[row],b[row]=t;
          for (j=k+1; j< n; j++){
               a[k][j]/=maxp;
               for (i=k+1;i<n;i++)
                   a[i][j] = a[i][k] * a[k][j];
          }
          b[k]/=maxp;
          for (i=k+1;i<n;i++)
              b[i]=b[k]*a[i][k];
     for (i=n-1;i>=0;i--)
          for (j=i+1;j< n;j++)
               b[i]=a[i][j]*b[j];
     for (k=0; k< n; k++)
         a[0][index[k]]=b[k];
     for (k=0; k< n; k++)
          b[k]=a[0][k];
     return 1;
```

# 14.5 线性相关

//判线性相关(正交化)

```
//传入 m 个 n 维向量
#include <math.h>
#define MAXN 100
#define eps 1e-10
int linear_dependent(int m,int n,double vec[][MAXN]){
     double ort[MAXN][MAXN],e;
     int i,j,k;
    if (m>n)
          return 1;
     for (i=0;i<m;i++){
          for (j=0; j< n; j++)
               ort[i][j]=vec[i][j];
          for (k=0; k< i; k++){
               for (e=j=0; j< n; j++)
                   e+=ort[i][j]*ort[k][j];
               for (j=0; j< n; j++)
                   ort[i][j]-=e*ort[k][j];
               for (e=j=0; j< n; j++)
                   e+=ort[i][j]*ort[i][j];
               if (fabs(e=sqrt(e))<eps)
                   return 1;
               for (j=0;j< n;j++)
                   ort[i][j]/=e;
          }
     }
     return 0;
}
14.6 日期
//日期函数
int days[12]=\{31,28,31,30,31,30,31,30,31,30,31\};
struct date {
     int year, month, day;
};
//判闰年
inline int leap(int year){
```

```
return (year%4==0&&year%100!=0)||year%400==0;
}
//判合法性
inline int legal(date a){
    if (a.month<0||a.month>12)
         return 0;
    if (a.month==2)
         return a.day>0&&a.day<=28+leap(a.year);
    return a.day>0&&a.day<=days[a.month-1];
}
//比较日期大小
inline int datecmp(date a,date b){
   if (a.year!=b.year)
        return a.year-b.year;
   if (a.month!=b.month)
        return a.month-b.month;
   return a.day-b.day;
}
//返回指定日期是星期几
int weekday(date a){
    int tm=a.month \ge 3?(a.month-2):(a.month+10);
    int ty=a.month>=3?a.year:(a.year-1);
    return (ty+ty/4-ty/100+ty/400+(int)(2.6*tm-0.2)+a.day)%7;
}
//日期转天数偏移
int date2int(date a){
    int ret=a.year*365+(a.year-1)/4-(a.year-1)/100+(a.year-1)/400,i;
    days[1] += leap(a.year);
    for (i=0;i\leq a.month-1;ret+=days[i++]);
    days[1]=28;
    return ret+a.day;
}
//天数偏移转日期
date int2date(int a){
    date ret;
    ret.year=a/146097*400;
    for (a\%=146097;a>=365+leap(ret.year);a=365+leap(ret.year),ret.year++);
    days[1]+=leap(ret.year);
    for (ret.month=1;a>=days[ret.month-1];a-=days[ret.month-1],ret.month++);
```

```
days[1]=28;
ret.day=a+1;
return ret;
}
```