# Team Note of Deobureo Minkyu Party

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#### 럭스를 럭스답게 든든한 연습헬팟 더불어민규당

```
hyea: Prove by solving.

koosaga: Locality to the rescue!

alex9801: Isshoman Beenzino
```

# 1 Flows, Matching

# 1.1 Hopcroft-Karp Bipartite Matching

```
const int MAXN = 50005, MAXM = 50005;
vector<int> gph[MAXN];
int dis[MAXN], 1[MAXN], r[MAXM], vis[MAXN];
void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
void add_edge(int 1, int r){ gph[1].push_back(r); }
bool bfs(int n){
  queue<int> que;
  bool ok = 0;
  memset(dis, 0, sizeof(dis));
  for(int i=0; i<n; i++){
   if(l[i] == -1 && !dis[i]){
      que.push(i);
      dis[i] = 1;
    }
  while(!que.empty()){
    int x = que.front();
    que.pop();
    for(auto &i : gph[x]){
      if(r[i] == -1) ok = 1;
      else if(!dis[r[i]]){
        dis[r[i]] = dis[x] + 1;
        que.push(r[i]);
  }
  return ok;
bool dfs(int x){
```

```
for(auto &i : gph[x]){
    if(r[i] == -1 \mid | (!vis[r[i]] \&\& dis[r[i]] == dis[x] + 1 \&\&
    dfs(r[i]))){
     vis[r[i]] = 1; l[x] = i; r[i] = x;
     return 1;
 }
 return 0;
int match(int n){
 memset(1, -1, sizeof(1));
 memset(r, -1, sizeof(r));
 int ret = 0;
 while(bfs(n)){
   memset(vis, 0, sizeof(vis));
   for(int i=0; i<n; i++) if(l[i] == -1 && dfs(i)) ret++;
 }
 return ret;
bool chk[MAXN + MAXM];
void rdfs(int x, int n){
 if(chk[x]) return;
 chk[x] = 1;
 for(auto &i : gph[x]){
    chk[i + n] = 1;
   rdfs(r[i], n);
vector<int> getcover(int n, int m){ // solve min. vertex cover
 match(n);
 memset(chk, 0, sizeof(chk));
 for(int i=0; i<n; i++) if(l[i] == -1) rdfs(i, n);
 vector<int> v;
 for(int i=0; i<n; i++) if(!chk[i]) v.push_back(i);</pre>
 for(int i=n; i<n+m; i++) if(chk[i]) v.push_back(i);</pre>
 return v:
```

# 1.2 Dinic's Algorithm

```
const int MAXN = 505;
struct edg{int pos, cap, rev;};
vector<edg> gph[MAXN];
void clear(){for(int i=0; i<MAXN; i++) gph[i].clear();}</pre>
void add_edge(int s, int e, int x){
  gph[s].push_back({e, x, (int)gph[e].size()});
  gph[e].push_back({s, 0, (int)gph[s].size()-1});
int dis[MAXN], pnt[MAXN];
bool bfs(int src, int sink){
  memset(dis, 0, sizeof(dis));
  memset(pnt, 0, sizeof(pnt));
  queue<int> que;
  que.push(src);
  dis[src] = 1;
  while(!que.empty()){
    int x = que.front();
    que.pop();
    for(auto &e : gph[x]){
      if(e.cap > 0 && !dis[e.pos]){
        dis[e.pos] = dis[x] + 1;
        que.push(e.pos);
      }
   }
  return dis[sink] > 0;
}
int dfs(int x, int sink, int f){
  if(x == sink) return f;
  for(; pnt[x] < gph[x].size(); pnt[x]++){</pre>
    edg e = gph[x][pnt[x]];
    if(e.cap > 0 \&\& dis[e.pos] == dis[x] + 1){
      int w = dfs(e.pos, sink, min(f, e.cap));
      if(w){
        gph[x][pnt[x]].cap -= w;
        gph[e.pos][e.rev].cap += w;
        return w;
```

```
return 0;
lint match(int src, int sink){
 lint ret = 0;
  while(bfs(src, sink)){
    int r;
    while((r = dfs(src, sink, 2e9))) ret += r;
 return ret;
     Min Cost Max Flow
const int MAXN = 100;
struct mincostflow{
  struct edg{ int pos, cap, rev, cost; };
  vector<edg> gph[MAXN];
 void clear(){
    for(int i=0; i<MAXN; i++) gph[i].clear();</pre>
  void add_edge(int s, int e, int x, int c){
    gph[s].push_back({e, x, (int)gph[e].size(), c});
    gph[e].push_back({s, 0, (int)gph[s].size()-1, -c});
 int dist[MAXN], pa[MAXN], pe[MAXN];
  bool inque[MAXN];
  bool spfa(int src, int sink){
   memset(dist, 0x3f, sizeof(dist));
    memset(inque, 0, sizeof(inque));
    queue<int> que;
    dist[src] = 0;
    inque[src] = 1;
    que.push(src);
    bool ok = 0;
    while(!que.empty()){
     int x = que.front();
      que.pop();
      if(x == sink) ok = 1;
```

```
inque[x] = 0:
      for(int i=0; i<gph[x].size(); i++){</pre>
        edg e = gph[x][i];
        if(e.cap > 0 \&\& dist[e.pos] > dist[x] + e.cost){
          dist[e.pos] = dist[x] + e.cost;
          pa[e.pos] = x;
          pe[e.pos] = i;
          if(!inque[e.pos]){
            inque[e.pos] = 1;
            que.push(e.pos);
         }
       }
    return ok;
  int match(int src, int sink){
    int ret = 0;
   while(spfa(src, sink)){
      int cap = 1e9;
      for(int pos = sink; pos != src; pos = pa[pos]){
        cap = min(cap, gph[pa[pos]][pe[pos]].cap);
      ret += dist[sink] * cap;
      for(int pos = sink; pos != src; pos = pa[pos]){
        int rev = gph[pa[pos]][pe[pos]].rev;
        gph[pa[pos]][pe[pos]].cap -= cap;
        gph[pos][rev].cap += cap;
    }
    return ret;
  }
}mcmf;
1.4 Hell-Joseon style MCMF
const int MAXN = 100;
struct mincostflow{
  struct edg{ int pos, cap, rev, cost; };
  vector<edg> gph[MAXN];
```

```
void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
void add_edge(int s, int e, int x, int c){
  gph[s].push_back({e, x, (int)gph[e].size(), c});
  gph[e].push_back({s, 0, (int)gph[s].size()-1, -c});
int phi[MAXN], inque[MAXN], dist[MAXN];
void prep(int src, int sink){
  memset(phi, 0x3f, sizeof(phi));
  memset(dist, 0x3f, sizeof(dist));
  queue<int> que;
  que.push(src);
  inque[src] = 1;
  while(!que.empty()){
    int x = que.front();
    que.pop();
    inque[x] = 0;
    for(auto &i : gph[x]){
      if(i.cap > 0 && phi[i.pos] > phi[x] + i.cost){
        phi[i.pos] = phi[x] + i.cost;
       if(!inque[i.pos]){
          inque[i.pos] = 1;
          que.push(i.pos);
  for(int i=0; i<MAXN; i++){</pre>
    for(auto &j : gph[i]){
      if(j.cap > 0) j.cost += phi[i] - phi[j.pos];
   }
  priority_queue<pi, vector<pi>, greater<pi> > pq;
  pq.push(pi(0, src));
  dist[src] = 0;
  while(!pq.empty()){
    auto 1 = pq.top();
    pq.pop();
    if(dist[1.second] != 1.first) continue;
    for(auto &i : gph[l.second]){
      if(i.cap > 0 && dist[i.pos] > 1.first + i.cost){
```

```
dist[i.pos] = 1.first + i.cost;
       pq.push(pi(dist[i.pos], i.pos));
     }
 }
bool vis[MAXN];
int ptr[MAXN];
int dfs(int pos, int sink, int flow){
 vis[pos] = 1;
 if(pos == sink) return flow;
 for(; ptr[pos] < gph[pos].size(); ptr[pos]++){</pre>
    auto &i = gph[pos][ptr[pos]];
    if(!vis[i.pos] && dist[i.pos] == i.cost + dist[pos] && i.cap >
    ) {
      int ret = dfs(i.pos, sink, min(i.cap, flow));
     if(ret != 0){
        i.cap -= ret;
        gph[i.pos][i.rev].cap += ret;
       return ret;
     }
   }
 }
 return 0;
int match(int src, int sink, int sz){
 prep(src, sink);
 for(int i=0; i<sz; i++) dist[i] += phi[sink] - phi[src];</pre>
 int ret = 0;
  while(true){
    memset(ptr, 0, sizeof(ptr));
    memset(vis, 0, sizeof(vis));
    int tmp = 0;
    while((tmp = dfs(src, sink, 1e9))){
     ret += dist[sink] * tmp;
     memset(vis, 0, sizeof(vis));
    tmp = 1e9;
    for(int i=0; i<sz; i++){</pre>
      if(!vis[i]) continue:
```

```
for(auto &j : gph[i]){
    if(j.cap > 0 && !vis[j.pos]){
        tmp = min(tmp, (dist[i] + j.cost) - dist[j.pos]);
    }
    }
    if(tmp > 1e9 - 200) break;
    for(int i=0; i<sz; i++){
        if(!vis[i]) dist[i] += tmp;
    }
} return ret;
}</pre>
```

#### 1.5 Circulation Problem

```
struct circ{
  maxflow mf;
  lint lsum;
  void clear(){
    lsum = 0:
    mf.clear();
  }
  void add_edge(int s, int e, int l, int r){
    lsum += 1:
    mf.add_edge(s + 2, e + 2, r - 1);
    mf.add\_edge(0, e + 2, 1);
    mf.add\_edge(s + 2, 1, 1);
  }
  bool solve(int s, int e){
    mf.add_edge(e+2, s+2, 1e9); // to reduce as maxflow with lower
    bounds, in circulation problem skip this line
    return lsum == mf.match(0, 1);
    // to get maximum LR flow, run maxflow from s+2 to e+2 again
  }
}circ;
```

# 1.6 Min Cost Circulation (WIP)

Should be added.

# 1.7 Gomory-Hu Tree

```
struct edg{ int s, e, x; };
vector<edg> edgs;
maxflow mf;
void clear(){edgs.clear();}
void add_edge(int s, int e, int x){edgs.push_back({s, e, x});}
bool vis[MAXN];
void dfs(int x){
  if(vis[x]) return;
  vis[x] = 1;
 for(auto &i : mf.gph[x]) if(i.cap > 0) dfs(i.pos);
vector<pi> solve(int n){
 // i - j cut : i - j minimum edge cost. 0 based.
  vector<pi> ret(n); // if i!=0, stores pair(parent,cost)
  for(int i=1; i<n; i++){
   for(auto &j : edgs){
      mf.add_edge(j.s, j.e, j.x);
      mf.add_edge(j.e, j.s, j.x);
    ret[i].first = mf.match(i, ret[i].second);
    memset(vis, 0, sizeof(vis));
    dfs(i):
    for(int j=i+1; j<n; j++){
      if(ret[j].second == ret[i].second && vis[j]){
        ret[i].second = i;
    mf.clear();
  return ret;
```

# 1.8 Stable Marriage

```
vector<vector<int>> rev; // 0-based
vector<int> solve(int n, vector<vector<int>> &a, vector<vector<int>>>
&b){
  rev.clear(); rev.resize(n, vector<int>(n, 0));
  vector<int> ret(n), pnt(n), mat(n, -1);
  queue<int> que;
 for(int i=0; i<n; i++){
   for(int j=0; j<n; j++) rev[i][b[i][j]] = j;</pre>
    que.push(i);
 }
  while(!que.empty()){
    int x = que.front(); que.pop();
    int y = a[x][pnt[x]++];
    if(mat[y] == -1) mat[y] = x, ret[x] = y;
      if(rev[y][mat[y]] > rev[y][x]){
        que.push(mat[y]);
       mat[y] = x, ret[x] = y;
      else que.push(x);
 return ret; // optimal matching of A side (x)
```

# 1.9 Edmond's Blossom Algorithm for General Matching (WIP)

Should be added.

# 2 Graph

# 2.1 2-SAT

```
strongly_connected scc;
int n; // = number of clauses
void init(int _n){ scc.clear(); n = _n; }
int NOT(int x){ return x >= n ? (x - n) : (x + n); }
```

```
void add_edge(int x, int y){ // input ~x to denote NOT
  if((x >> 31) & 1) x = (x) + n;
 if((y >> 31) & 1) y = (~y) + n;
  scc.add_edge(x, y), scc.add_edge(NOT(y), NOT(x));
bool satisfy(vector<bool> &res){
  res.resize(n);
  scc.get_scc(2*n);
  for(int i=0; i<n; i++){</pre>
   if(scc.comp[i] == scc.comp[NOT(i)]) return 0;
   if(scc.comp[i] < scc.comp[NOT(i)]) res[i] = 0;</pre>
    else res[i] = 1;
  }
 return 1;
}
2.2
     BCC
void color(int x, int p){
  if(p){
    bcc[p].push_back(x);
    cmp[x].push_back(p);
  for(auto &i : gph[x]){
    if(cmp[i].size()) continue;
    if(low[i] >= dfn[x]){
      bcc[++c].push_back(x);
      cmp[x].push_back(c);
      color(i, c);
    else color(i, p);
 }
}
     Splay Tree + Link-Cut Tree
// Checklist 1. Is it link cut, or splay?
// Checklist 2. In link cut, is son always root?
void rotate(node *x){
  if(!x->p) return;
```

```
push(x->p); // if there's lazy stuff
  push(x);
  node *p = x->p;
  bool is_left = (p->l == x);
  node *b = (is_left ? x->r : x->l);
  x->p = p->p;
  if(x-p \&\& x-p->l == p) x-p->l = x;
  if(x-p \&\& x-p-r == p) x-p-r = x;
  if(is_left){
   if(b) b->p = p;
   p->1 = b;
   p->p = x;
    x->r = p;
  else{
    if(b) b \rightarrow p = p;
   p->r = b;
   p->p = x;
    x->1 = p;
  pull(p); // if there's something to pull up
  pull(x);
 if(!x->p) root = x; // IF YOU ARE SPLAY TREE
 if(p->pp){ // IF YOU ARE LINK CUT TREE
    x->pp = p->pp;
   p->pp = NULL;
void splay(node *x){
  while(x->p){
    node *p = x->p;
    node *g = p - p;
    if(g){
     if((p->1 == x) ^ (g->1 == p)) rotate(x);
      else rotate(p);
    rotate(x);
 }
void access(node *x){
```

```
splay(x);
  push(x);
  if(x->r){
    x->r->pp = x;
    x->r->p = NULL;
    x->r = NULL;
  pull(x);
  while(x->pp){
    node *nxt = x->pp;
    splay(nxt);
    push(nxt);
    if(nxt->r){
      nxt->r->pp = nxt;
      nxt->r->p = NULL;
      nxt->r = NULL;
    }
    nxt->r = x;
    x->p = nxt;
    x->pp = NULL;
    pull(nxt);
    splay(x);
}
node *root(node *x){
  access(x);
  while (x->1) {
    push(x);
    x = x->1;
  access(x);
  return x;
node *par(node *x){
  access(x);
  if(!x->1) return NULL;
 push(x);
  x = x->1;
  while (x->r) {
    push(x);
```

```
x = x->r:
 access(x);
 return x;
node *lca(node *s, node *t){
 access(s);
 access(t);
  splay(s);
 if(s->pp == NULL) return s;
 return s->pp;
void link(node *par, node *son){
 access(par);
  access(son);
  son->rev ^= 1; // remove if needed
 push(son);
 son->1 = par;
 par->p = son;
 pull(son);
void cut(node *p){
 access(p);
 push(p);
 if(p->1){
   p->1->p = NULL;
   p->1 = NULL;
 pull(p);
```

# 2.4 Offline Dynamic MST (WIP)

Should be added.

# 2.5 Dominator Tree

```
namespace dtree{    // by cki86201
    vector<int> E[MAXN], RE[MAXN], rdom[MAXN];
```

```
int S[MAXN], RS[MAXN], cs;
int par[MAXN], val[MAXN], sdom[MAXN], rp[MAXN], dom[MAXN];
void clear(int n) {
 cs = 0;
 for(int i=0;i<=n;i++) {</pre>
    par[i] = val[i] = sdom[i] = rp[i] = dom[i] = S[i] = RS[i] = 0;
   E[i].clear(); RE[i].clear(); rdom[i].clear();
 }
void add_edge(int x, int y) { E[x].push_back(y); }
void Union(int x, int y) { par[x] = y; }
int Find(int x, int c = 0) {
  if(par[x] == x) return c ? -1 : x;
 int p = Find(par[x], 1);
 if (p == -1) return c ? par[x] : val[x];
 if(sdom[val[x]] > sdom[val[par[x]]]) val[x] = val[par[x]];
 par[x] = p;
 return c ? p : val[x];
void dfs(int x) {
 RS[S[x] = ++cs] = x:
 par[cs] = sdom[cs] = val[cs] = cs;
 for(int e : E[x]) {
   if(S[e] == 0) dfs(e), rp[S[e]] = S[x];
    RE[S[e]].push_back(S[x]);
 }
}
int solve(int s, int *up) { // Calculate idoms
  dfs(s);
 for(int i=cs;i;i--) {
    for(int e : RE[i]) sdom[i] = min(sdom[i], sdom[Find(e)]);
    if(i > 1) rdom[sdom[i]].push_back(i);
    for(int e : rdom[i]) {
      int p = Find(e);
     if(sdom[p] == i) dom[e] = i;
      else dom[e] = p;
    if(i > 1) Union(i, rp[i]);
 }
```

```
for(int i=2;i<=cs;i++) if(sdom[i] != dom[i]) dom[i] =</pre>
    dom[dom[i]]:
   for(int i=2;i<=cs;i++) up[RS[i]] = RS[dom[i]];</pre>
   return cs;
 }
     Global Min-Cut
namespace stoer_wagner{
 int minimum_cut_phase(int n, int &s, int &t, vector<vector<int>>
 &adj, vector<int> vis){
    vector<int> dist(n);
   int mincut = 1e9;
    while(true){
     int pos = -1, cur = -1e9;
     for(int i=0; i<n; i++){
       if(!vis[i] && dist[i] > cur){
          cur = dist[i];
         pos = i;
     if(pos == -1) break;
     s = t:
     t = pos;
     mincut = cur;
     vis[pos] = 1;
     for(int i=0; i<n; i++){
       if(!vis[i]) dist[i] += adj[pos][i];
     }
    return mincut; // optimal s-t cut here is, {t} and V \ {t}
 int solve(int n, vector<vector<int>> adj){
    if(n <= 1) return 0;
   vector<int> vis(n);
   int ans = 1e9;
    for(int i=0; i<n-1; i++){</pre>
```

ans = min(ans, minimum\_cut\_phase(n, s, t, adj, vis));

int s, t;

```
vis[t] = 1;
for(int j=0; j<n; j++){
    if(!vis[j]){
        adj[s][j] += adj[t][j];
        adj[j][s] += adj[j][t];
    }
    }
    adj[s][s] = 0;
}
return ans;
}
};</pre>
```

# 2.7 K-shortest path (WIP)

Should be added.

# 2.8 Edmond's Directed MST (WIP)

Should be added.

# 2.9 Vizing Theorem for Edge Coloring (WIP)

Should be added.

# 3 Strings

# 3.1 Aho-Corasick Algorithm

```
const int MAXN = 100005, MAXC = 26;
int trie[MAXN][MAXC], fail[MAXN], term[MAXN], piv;
void init(vector<string> &v){
  memset(trie, 0, sizeof(trie));
  memset(fail, 0, sizeof(fail));
  memset(term, 0, sizeof(term));
  piv = 0;
  for(auto &i : v){
    int p = 0;
    for(auto &j : i){
        if(!trie[p][j]) trie[p][j] = ++piv;
    }
}
```

```
p = trie[p][j];
    term[p] = 1;
  queue<int> que;
  for(int i=0; i<MAXC; i++){</pre>
    if(trie[0][i]) que.push(trie[0][i]);
  while(!que.empty()){
    int x = que.front();
    que.pop();
    for(int i=0; i<MAXC; i++){</pre>
      if(trie[x][i]){
        int p = fail[x];
        while(p && !trie[p][i]) p = fail[p];
        p = trie[p][i];
        fail[trie[x][i]] = p;
        if(term[p]) term[trie[x][i]] = 1;
        que.push(trie[x][i]);
bool query(string &s){
 int p = 0;
 for(auto &i : s){
    while(p && !trie[p][i]) p = fail[p];
    p = trie[p][i];
    if(term[p]) return 1;
 return 0;
    Suffix Array
 Should be revised.
const int MAXN = 500005;
int ord[MAXN], nord[MAXN], cnt[MAXN], aux[MAXN];
void solve(int n, char *str, int *sfx, int *rev, int *lcp){
 int p = 1;
```

```
memset(ord, 0, sizeof(ord));
for(int i=0; i<n; i++){
  sfx[i] = i:
  ord[i] = str[i];
int pnt = 1;
while(1){
  memset(cnt, 0, sizeof(cnt));
  for(int i=0; i<n; i++) cnt[ord[min(i+p, n)]]++;</pre>
  for(int i=1; i<=n || i<=255; i++) cnt[i] += cnt[i-1];
  for(int i=n-1; i>=0; i--)
    aux[--cnt[ord[min(i+p, n)]]] = i;
  memset(cnt, 0, sizeof(cnt));
  for(int i=0; i<n; i++) cnt[ord[i]]++;</pre>
  for(int i=1; i<=n || i<=255; i++) cnt[i] += cnt[i-1];
  for(int i=n-1: i>=0: i--)
    sfx[--cnt[ord[aux[i]]]] = aux[i];
  if(pnt == n) break;
  pnt = 1;
  nord[sfx[0]] = 1;
  for(int i=1; i<n; i++){
    if(ord[sfx[i-1]] != ord[sfx[i]] || ord[sfx[i-1] + p] !=
    ord[sfx[i] + p]){
     pnt++;
    nord[sfx[i]] = pnt;
  memcpy(ord, nord, sizeof(int) * n);
 p *= 2;
for(int i=0; i<n; i++) rev[sfx[i]] = i;</pre>
int h = 0:
for(int i=0; i<n; i++){
  if(rev[i]){
    int prv = sfx[rev[i] - 1];
    while(str[prv + h] == str[i + h]) h++;
    lcp[rev[i]] = h;
 h = \max(h-1, 0);
```

# 3.3 Manacher's Algorithm

```
const int MAXN = 1000005;
int aux[2 * MAXN - 1];
void solve(int n, int *str, int *ret){
 // *ret : number of nonobvious palindromic character pair
 for(int i=0; i<n; i++){
    aux[2*i] = str[i];
   if(i != n-1) aux[2*i+1] = -1;
 }
  int p = 0, c = 0;
 for(int i=0; i<2*n-1; i++){
    int cur = 0:
    if(i \le p) cur = min(ret[2 * c - i], p - i);
    while(i - cur - 1 >= 0 && i + cur + 1 < 2*n-1 && aux[i-cur-1] ==
    aux[i+cur+1]){
      cur++;
    ret[i] = cur;
    if(i + ret[i] > p){
     p = i + ret[i];
     c = i;
```

# 3.4 Suffix Automaton (WIP)

Should be **added**.

# 3.5 Palindromic Tree (WIP)

Should be **added**.

# 3.6 Circular LCS

```
string s1, s2;
int dp[4005][2005];
```

```
int nxt[4005][2005];
int n, m;
void reroot(int px){
  int py = 1;
  while(py <= m && nxt[px][py] != 2) py++;
  nxt[px][py] = 1;
  while(px < 2 * n \&\& py < m){
    if(nxt[px+1][py] == 3){
      px++;
      nxt[px][py] = 1;
    else if(nxt[px+1][py+1] == 2){
      px++;
      py++;
      nxt[px][py] = 1;
    else py++;
  while(px < 2 * n && nxt[px+1][py] == 3){
    px++;
    nxt[px][py] = 1;
}
int track(int x, int y, int e){ // use this routine to find LCS as
string
  int ret = 0;
  while(v != 0 \&\& x != e){
    if(nxt[x][y] == 1) y--;
   else if(nxt[x][y] == 2) ret += (s1[x] == s2[y]), x--, y--;
    else if(nxt[x][y] == 3) x--;
  }
  return ret;
}
int solve(string a, string b){
  n = a.size(), m = b.size();
  s1 = "#" + a + a;
  s1 = '#' + b:
```

```
for(int i=0; i<=2*n; i++){
  for(int j=0; j<=m; j++){</pre>
    if(j == 0){
      nxt[i][j] = 3;
      continue;
    if(i == 0){
     nxt[i][j] = 1;
      continue;
    dp[i][j] = -1;
   if(dp[i][j] < dp[i][j-1]){</pre>
      dp[i][j] = dp[i][j-1];
      nxt[i][j] = 1;
    if(dp[i][j] < dp[i-1][j-1] + (s1[i] == s2[j])){
      dp[i][j] = dp[i-1][j-1] + (s1[i] == s2[j]);
      nxt[i][j] = 2;
    if(dp[i][j] < dp[i-1][j]){
      dp[i][j] = dp[i-1][j];
      nxt[i][j] = 3;
int ret = dp[n][m];
for(int i=1; i<n; i++){
  reroot(i), ret = max(ret, track(n+i, m, i));
return ret;
```

# 4 Geometry

# 4.1 Green's Theorem

Let C is positive, smooth, simple curve. D is region bounded by C.

$$\oint_C (Ldx + Mdy) = \iint_D (\frac{\partial M}{\partial x} - \frac{\partial L}{\partial y})$$

}

Circle solve(vector<Point> p) {

Circle c = INVAL;

```
To calculate area, \frac{\partial M}{\partial x} - \frac{\partial L}{\partial y} = 1, common selection is M = \frac{1}{2}x, L = -\frac{1}{2}y.
  Line integral of circle parametrized by (x, y) = (x_C + r_C \cos \theta, y_C + r_C \sin \theta) is
given as follows: \frac{1}{2}(r_C(x_C(\sin\theta_f - \sin\theta_i) - y_C(\cos\theta_f - \cos\theta_i)) + (\theta_f - \theta_i)r_C^2).
  Line integral of line parametrized by (x,y) = t(x_1,y_1) + (1-t)(x_2,y_2) is given
as follows.: \frac{1}{2}(x_1y_2 - x_2y_1).
inline double arc_area(double x, double y, double r, double s,
double e){ //s and e are line integral theta value.
  return (r*(x*(sin(e)-sin(s))-v*(cos(e)-cos(s)))+(e-s)*r*r)*0.5:
}
inline double polygon_area(double x1, double y1, double x2, double
y2){ //Shoelace formula
  return (x1*y2-x2*y1)*0.5;
}
      Smallest Enclosing Circle / Sphere
namespace cover_2d{
    double eps = 1e-9;
    using Point = complex<double>;
    struct Circle{ Point p; double r; };
    double dist(Point p, Point q){ return abs(p-q); }
    double area2(Point p, Point q){ return (conj(p)*q).imag();}
    bool in(const Circle& c, Point p){ return dist(c.p, p) < c.r +</pre>
    eps: }
    Circle INVAL = Circle{Point(0, 0), -1};
    Circle mCC(Point a, Point b, Point c){
         b -= a; c -= a;
         double d = 2*(conj(b)*c).imag(); if(abs(d) < eps) return
         INVAL;
         Point ans = (c*norm(b) - b*norm(c)) * Point(0, -1) / d;
         return Circle{a + ans, abs(ans)};
```

mt19937 gen(0x94949); shuffle(p.begin(), p.end(), gen);

```
for(int i=0; i<p.size(); ++i) if(c.r<0 ||!in(c, p[i])){
            c = Circle{p[i]. 0}:
            for(int j=0; j<=i; ++j) if(!in(c, p[j])){</pre>
                Circle ans{(p[i]+p[j])*0.5, dist(p[i], p[j])*0.5};
                if(c.r == 0) {c = ans; continue;}
                Circle 1, r; 1 = r = INVAL;
                Point pq = p[j]-p[i];
                for(int k=0; k<=j; ++k) if(!in(ans, p[k])) {</pre>
                     double a2 = area2(pq, p[k]-p[i]);
                     Circle c = mCC(p[i], p[j], p[k]);
                     if(c.r<0) continue;</pre>
                     else if(a2 > 0 && (1.r<0||area2(pg, c.p-p[i]) >
                     area2(pq, l.p-p[i]))) l = c;
                     else if(a2 < 0 && (r.r<0||area2(pg, c.p-p[i]) <
                     area2(pq, r.p-p[i])) r = c;
                if(1.r<0&&r.r<0) c = ans:
                else if(1.r<0) c = r;
                else if(r.r<0) c = 1;
                else c = 1.r<=r.r?1:r;</pre>
            }
        }
        return c;
};
namespace cover_3d{
  double enclosing_sphere(vector<double> x, vector<double> y,
  vector<double> z){
    int n = x.size();
    auto hyp = [](double x, double y, double z){
      return x * x + y * y + z * z;
    };
    double px = 0, py = 0, pz = 0;
    for(int i=0: i<n: i++){</pre>
      px += x[i]:
      py += y[i];
      pz += z[i];
    px *= 1.0 / n:
```

int a, b;

```
pv *= 1.0 / n:
    pz *= 1.0 / n;
    double rat = 0.1, maxv;
    for(int i=0; i<10000; i++){
      maxv = -1;
      int maxp = -1;
      for(int j=0; j<n; j++){
        double tmp = hyp(x[j] - px, y[j] - py, z[j] - pz);
        if(maxv < tmp){</pre>
          maxv = tmp;
          maxp = j;
      px += (x[maxp] - px) * rat;
      py += (y[maxp] - py) * rat;
      pz += (z[maxp] - pz) * rat;
      rat *= 0.998:
    return sqrt(maxv);
};
     3D Convex Hull
// code credit : https://gist.github.com/msg555/4963794
struct vec3{
  11 x, y, z;
  vec3(): x(0), y(0), z(0) {}
  vec3(11 a, 11 b, 11 c): x(a), y(b), z(c) {}
  vec3 operator*(const vec3& v) const{ return vec3(y*v.z-z*v.y,
  z*v.x-x*v.z, x*v.y-y*v.x); }
  vec3 operator-(const vec3& v) const{ return vec3(x-v.x, y-v.y,
  z-v.z): }
  vec3 operator-() const{ return vec3(-x, -v, -z); }
  11 dot(const vec3 &v) const{ return x*v.x+y*v.y+z*v.z; }
};
struct twoset {
```

void insert(int x) {  $(a == -1 ? a : b) = x; }$ 

```
bool contains(int x) { return a == x || b == x; }
  void erase(int x) { (a == x ? a : b) = -1: }
 int size() { return (a != -1) + (b != -1); }
} E[MAXN] [MAXN]; // i < j</pre>
struct face{
  vec3 norm;
 ll disc;
 int I[3];
};
face make_face(int i, int j, int k, int ii, vector<vec3> &A){ // p^T
* norm < disc
  E[i][j].insert(k); E[i][k].insert(j); E[j][k].insert(i);
 face f; f.I[0]=i, f.I[1]=j, f.I[2]=k;
 f.norm = (A[j]-A[i])*(A[k]-A[i]);
 f.disc = f.norm.dot(A[i]);
 if(f.norm.dot(A[ii])>f.disc){
   f.norm = -f.norm;
   f.disc = -f.disc;
 return f:
vector<face> get_hull(vector<vec3> &A){
 int N = A.size();
  vector<face> faces; memset(E, -1, sizeof(E));
  faces.push_back(make_face(0,1,2,3,A));
  faces.push_back(make_face(0,1,3,2,A));
  faces.push_back(make_face(0,2,3,1,A));
  faces.push_back(make_face(1,2,3,0,A));
  for(int i=4; i<N; ++i){
    for(int j=0; j<faces.size(); ++j){</pre>
      face f = faces[i];
      if(f.norm.dot(A[i])>f.disc){
        E[f.I[0]][f.I[1]].erase(f.I[2]):
        E[f.I[0]][f.I[2]].erase(f.I[1]);
        E[f.I[1]][f.I[2]].erase(f.I[0]);
        faces[j--] = faces.back();
        faces.pop_back();
```

```
}
  }
 int nf = faces.size();
 for(int j=0; j<nf; ++j){</pre>
   face f=faces[j];
   for(int a=0; a<3; ++a) for(int b=a+1; b<3; ++b){
      int c=3-a-b;
      if(E[f.I[a]][f.I[b]].size()==2) continue;
     faces.push_back(make_face(f.I[a], f.I[b], i, f.I[c], A));
 }
return faces;
  Dynamic Convex Hull Trick
```

```
// code credit : https://github.com/niklasb/contest-algos/
// blob/master/convex_hull/dynamic.cpp
using line_t = double;
const line_t is_query = -1e18;
struct Line {
  line_t m, b;
  mutable function<const Line*()> succ;
  bool operator<(const Line& rhs) const {</pre>
    if (rhs.b != is_query) return m < rhs.m;</pre>
    const Line* s = succ();
    if (!s) return 0:
    line_t x = rhs.m;
    return b - s -> b < (s -> m - m) * x:
 }
};
struct HullDynamic : public multiset<Line> { // will maintain upper
hull for maximum
  bool bad(iterator y) {
    auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
```

```
return y->m == z->m && y->b <= z->b;
    auto x = prev(y);
    if (z == end()) return y->m == x->m && y->b <= x->b;
   return (x-b - y-b)*(z-m - y-m) >= (y-b - z-b)*(y-m - y-m)
   x->m);
 }
 void insert_line(line_t m, line_t b) {
    auto y = insert({ m, b });
   y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
    if (bad(y)) { erase(y); return; }
    while (next(y) != end() && bad(next(y))) erase(next(y));
    while (y != begin() && bad(prev(y))) erase(prev(y));
 line_t query(line_t x) {
    auto 1 = *lower_bound((Line) { x, is_query });
   return 1.m * x + 1.b;
 }
}H;
```

# 4.5 Half-plane Intersection

```
const double eps = 1e-8;
typedef pair<long double, long double> pi;
namespace hpi{
  bool z(long double x){ return fabs(x) < eps; }</pre>
  struct line{
    long double a, b, c;
    bool operator<(const line &1)const{</pre>
      bool flag1 = pi(a, b) > pi(0, 0);
      bool flag2 = pi(1.a, 1.b) > pi(0, 0);
      if(flag1 != flag2) return flag1 > flag2;
      long double t = ccw(pi(0, 0), pi(a, b), pi(1.a, 1.b));
      return z(t) ? c * hypot(1.a, 1.b) < 1.c * hypot(a, b) : t > 0;
    pi slope(){ return pi(a, b);}
  pi cross(line a, line b){
    long double det = a.a * b.b - b.a * a.b;
```

```
return pi((a.c * b.b - a.b * b.c) / det, (a.a * b.c - a.c * b.a) 4.7 kd-tree
    / det):
  }
  bool bad(line a, line b, line c){
    if(ccw(pi(0, 0), a.slope(), b.slope()) <= 0) return false;</pre>
    pi crs = cross(a, b);
    return crs.first * c.a + crs.second * c.b >= c.c;
  bool solve(vector<line> v, vector<pi> &solution){ // ax + by <= c;</pre>
    sort(v.begin(), v.end());
    deque<line> dq;
   for(auto &i : v){
      if(!dq.empty() && z(ccw(pi(0, 0), dq.back().slope(),
      i.slope()))) continue;
      while(dg.size() >= 2 && bad(dg[dg.size()-2], dg.back(), i))
      dq.pop_back();
      while(dq.size() >= 2 \&\& bad(i, dq[0], dq[1])) dq.pop_front();
      dq.push_back(i);
    while(dq.size() > 2 \&\& bad(dq[dq.size()-2], dq.back(), dq[0]))
    dq.pop_back();
    while(dq.size() > 2 && bad(dq.back(), dq[0], dq[1]))
    dq.pop_front();
    vector<pi> tmp;
    for(int i=0; i<dq.size(); i++){</pre>
      line cur = dq[i], nxt = dq[(i+1)%dq.size()];
      if(ccw(pi(0, 0), cur.slope(), nxt.slope()) <= eps) return</pre>
      false:
      tmp.push_back(cross(cur, nxt));
    solution = tmp;
    return true;
};
    Polygon tangent (WIP)
```

Should be added.

```
typedef pair<int, int> pi;
struct node{
  pi pnt;
  int spl, sx, ex, sy, ey;
}tree[270000];
pi a[100005]:
int n. ok[270000]:
lint sqr(int x){ return 111 * x * x; }
bool cmp1(pi a, pi b){ return a < b; }</pre>
bool cmp2(pi a, pi b){ return pi(a.second, a.first) < pi(b.second,
b.first); }
// init(0, n-1, 1) : Initialize kd-tree
// set dap = INF, and call solve(1, P). dap = (closest point from P)
void init(int s, int e, int p){ // Initialize kd-tree
  int minx = 1e9, maxx = -1e9, miny = 1e9, maxy = -1e9;
  int m = (s+e)/2;
  for(int i=s; i<=e; i++){
    minx = min(minx, a[i].first);
    miny = min(miny, a[i].second);
    maxx = max(maxx, a[i].first);
    maxy = max(maxy, a[i].second);
  tree[p].spl = (maxx - minx < maxy - miny);</pre>
  sort(a+s, a+e+1, [&](const pi &a, const pi &b){
    return tree[p].spl ? cmp2(a, b) : cmp1(a, b);
  });
  ok[p] = 1;
  tree[p] = {a[m], tree[p].spl, minx, maxx, miny, maxy};
 if (s \le m-1) init(s, m-1, 2*p);
  if(m+1 \le e) init(m+1, e, 2*p+1);
lint dap = 3e18;
```

```
void solve(int p, pi x){ // find closest point from point x (L^2)
  if(x != tree[p].pnt) dap = min(dap, sqr(x.first -
  tree[p].pnt.first) + sqr(x.second - tree[p].pnt.second));
  if(tree[p].spl){
    if(!cmp2(tree[p].pnt, x)){
      if(ok[2*p]) solve(2*p, x);
      if (ok[2*p+1] \&\& sqr(tree[2*p+1].sy - x.second) < dap)
      solve(2*p+1, x);
    }
    else{
      if(ok[2*p+1]) solve(2*p+1, x);
      if (ok[2*p] \&\& sqr(tree[2*p].ey - x.second) < dap) solve(2*p,
      x);
    }
  }
  else{
    if(!cmp1(tree[p].pnt, x)){
      if(ok[2*p]) solve(2*p, x);
      if (ok[2*p+1] \&\& sqr(tree[2*p+1].sx - x.first) < dap)
      solve(2*p+1, x);
    }
    else{
      if(ok[2*p+1]) solve(2*p+1, x);
      if(ok[2*p] && sqr(tree[2*p].ex - x.first) < dap) solve(2*p,</pre>
      x);
    }
 }
}
5
    Math
5.1 FFT / NTT
namespace fft{
  typedef complex<double> base;
  void fft(vector<base> &a, bool inv){
    int n = a.size(), j = 0;
    vector<base> roots(n/2);
    for(int i=1; i<n; i++){</pre>
      int bit = (n >> 1);
```

```
while(j >= bit){
      j -= bit;
      bit >>= 1;
    j += bit;
   if(i < j) swap(a[i], a[j]);</pre>
  double ang = 2 * acos(-1) / n * (inv ? -1 : 1);
  for(int i=0; i<n/2; i++){
    roots[i] = base(cos(ang * i), sin(ang * i));
  /* In NTT, let prr = primitive root. Then,
  int ang = ipow(prr, (mod - 1) / n);
  if(inv) ang = ipow(ang, mod - 2);
  for(int i=0; i< n/2; i++){
   roots[i] = (i ? (111 * roots[i-1] * ang % mod) : 1);
  Others are same. If there is /= n, do *= ipow(n, mod - 2).
  In XOR convolution, roots[*] = 1.
  */
  for(int i=2; i<=n; i<<=1){</pre>
   int step = n / i;
   for(int j=0; j<n; j+=i){
     for(int k=0; k<i/2; k++){
        base u = a[j+k], v = a[j+k+i/2] * roots[step * k];
        a[j+k] = u+v;
        a[j+k+i/2] = u-v;
  if(inv) for(int i=0; i<n; i++) a[i] /= n;
vector<lint> multiply(vector<lint> &v, vector<lint> &w){
  vector<base> fv(v.begin(), v.end()), fw(w.begin(), w.end());
  int n = 2; while (n < v.size() + w.size()) n <<= 1;
  fv.resize(n); fw.resize(n);
  fft(fv, 0); fft(fw, 0);
  for(int i=0; i<n; i++) fv[i] *= fw[i];
  fft(fv. 1):
```

```
vector<lint> ret(n):
 for(int i=0: i<n: i++) ret[i] = (lint)round(fv[i].real()):</pre>
 return ret:
vector<lint> multiply(vector<lint> &v, vector<lint> &w, lint mod){
 int n = 2; while(n < v.size() + w.size()) n <<= 1;</pre>
 vector \langle base \rangle v1(n), v2(n), r1(n), r2(n);
 for(int i=0; i<v.size(); i++){</pre>
    v1[i] = base(v[i] >> 15, v[i] & 32767);
 for(int i=0; i<w.size(); i++){</pre>
    v2[i] = base(w[i] >> 15, w[i] & 32767);
 fft(v1, 0);
 fft(v2, 0);
 for(int i=0; i<n; i++){
   int j = (i ? (n - i) : i);
   base ans1 = (v1[i] + conj(v1[j])) * base(0.5, 0);
    base ans2 = (v1[i] - conj(v1[j])) * base(0, -0.5);
    base ans3 = (v2[i] + conj(v2[j])) * base(0.5, 0);
   base ans4 = (v2[i] - conj(v2[j])) * base(0, -0.5);
   r1[i] = (ans1 * ans3) + (ans1 * ans4) * base(0, 1);
    r2[i] = (ans2 * ans3) + (ans2 * ans4) * base(0, 1);
 fft(r1, 1);
 fft(r2, 1);
 vector<lint> ret(n);
 for(int i=0; i<n; i++){
   lint av = (lint)round(r1[i].real());
   lint bv = (lint)round(r1[i].imag()) +
    (lint)round(r2[i].real());
    lint cv = (lint)round(r2[i].imag());
    av %= mod, bv %= mod, cv %= mod;
    ret[i] = (av << 30) + (bv << 15) + cv;
    ret[i] %= mod:
   ret[i] += mod;
    ret[i] %= mod;
 return ret;
```

# 5.2 Hell-Joseon style FFT

```
#include <smmintrin.h>
#include <immintrin.h>
#pragma GCC target("avx2")
#pragma GCC target("fma")
__m256d mult(__m256d a, __m256d b){
    _{m256d} c = _{mm256_{movedup_{pd}(a)}}
    _{m256d} d = _{mm256\_shuffle\_pd(a, a, 15)};
    _{m256d} cb = _{mm256} ul_{pd}(c, b);
    _{m256d} db = _{mm256} dd, b);
    _{m256d} = _{mm256\_shuffle\_pd(db, db, 5)};
    _{\rm m256d} r = _{\rm mm256\_addsub\_pd(cb, e)};
    return r;
void fft(int n, __m128d a[], bool invert){
    for(int i=1, j=0; i<n; ++i){
        int bit = n >> 1;
        for(; j>=bit; bit>>=1) j -= bit;
        j += bit;
        if(i<j) swap(a[i], a[j]);
    for(int len=2; len<=n; len<<=1){</pre>
        double ang = 2*3.14159265358979/len*(invert?-1:1);
        _{\text{m256d wlen}}; wlen[0] = cos(ang), wlen[1] = sin(ang);
        for(int i=0; i<n; i += len){</pre>
            _{m256d} w; w[0] = 1; w[1] = 0;
            for(int j=0; j<len/2; ++j){
                 w = _mm256_permute2f128_pd(w, w, 0);
                 wlen = _{mm256}_{insertf128}_{pd(wlen, a[i+j+len/2], 1)};
                 w = mult(w, wlen);
                 _{m128d} vw = _{mm256} extractf128_{pd}(w, 1);
                 _{\rm m}128d\ u = a[i+j];
                 a[i+j] = _mm_add_pd(u, vw);
                 a[i+j+len/2] = _mm_sub_pd(u, vw);
            }
```

```
if(invert){
        _{m128d inv; inv[0] = inv[1] = 1.0/n;
        for(int i=0; i<n; ++i) a[i] = _mm_mul_pd(a[i], inv);</pre>
    }
}
vector<int64_t> multiply(vector<int64_t>& v, vector<int64_t>& w){
    int n = 2; while(n < v.size()+w.size()) n <<=1;</pre>
    _{m128d* fv = new _{m128d[n];}
    for(int i=0; i<n; ++i) fv[i][0] = fv[i][1] = 0;
    for(int i=0; i<v.size(); ++i) fv[i][0] = v[i];</pre>
    for(int i=0; i<w.size(); ++i) fv[i][1] = w[i];</pre>
    fft(n, fv, 0); //(a+bi) is stored in FFT
    for(int i=0; i<n; i += 2){
        __m256d a;
        a = _mm256_insertf128_pd(a, fv[i], 0);
        a = _{mm256}insertf128_pd(a, fv[i+1], 1);
        a = mult(a, a);
        fv[i] = _mm256_extractf128_pd(a, 0);
        fv[i+1] = _mm256_extractf128_pd(a, 1);
    }
    fft(n, fv, 1):
    vector<int64 t> ret(n):
    for(int i=0; i<n; ++i) ret[i] = (int64_t)round(fv[i][1]/2);
    delete[] fv;
    return ret;
}
     NTT Polynomial Division
```

```
vector<lint> get_inv(int n, const vector<lint> &p){
  vector<lint> q = {ipow(p[0], mod - 2)};
  for(int i=2; i<=n; i<<=1){
    vector<lint> res;
    vector<lint> fq(q.begin(), q.end()); fq.resize(2*i);
    vector<lint> fp(p.begin(), p.begin() + i); fp.resize(2*i);
  fft(fq, 0); fft(fp, 0);
  for(int j=0; j<2*i; j++){
    fp[j] *= fq[j] * fq[j] % mod;
    fp[j] %= mod;
}</pre>
```

```
fft(fp, 1);
    res.resize(i):
   for(int j=0; j<i; j++){
     res[j] = mod - fp[j];
     if(j < i/2) res[j] += 2 * q[j];
     res[j] %= mod;
    q = res;
 return q;
vector<lint> poly_divide(const vector<lint> &a, const vector<lint>
&b){
 assert(b.back() != 0); // please trim leading zero
 int n = a.size(), m = b.size();
 int k = 2; while (k < n-m+1) k <<= 1;
 vector<lint> rb(k), ra(k);
 for(int i=0; i<m && i<k; ++i) rb[i] = b[m-i-1];
 for(int i=0; i<n && i<k; ++i) ra[i] = a[n-i-1];
 vector<lint> rbi = get_inv(k, rb);
 vector<lint> res = multiplv(rbi, ra):
 res.resize(n - m + 1):
 reverse(res.begin(), res.end());
 return res;
5.4 Simplex Algorithm
/* Ax <= b. max c^T x
* Usage : Simplex(VVD A, VD b, VD c).solve(VD ans)
* not feasible : -INF; unbounded : INF
* accuracy ~ (size of ans) * EPS
* EPS recommended 1e-9 on double, 1e-12 on long double
```

```
/* Ax <= b, max c^T x
 * Usage : Simplex(VVD A, VD b, VD c).solve(VD ans)
 * not feasible : -INF; unbounded : INF
 * accuracy ~ (size of ans) * EPS
 * EPS recommended 1e-9 on double, 1e-12 on long double
 * expected n ~ 100, 10ms. worst case is exponential */
using real_t = double;
using VD = vector<real_t>;
using VVD = vector<VD>;
const real_t EPS = 1e-9;
```

```
struct Simplex{
  int m, n;
  vector<int> B, N;
  VVD D;
  Simplex(const VVD& A, const VD& b, const VD &c)
    : m(b.size()), n(c.size()), N(n+1), B(m), D(m+2, VD(n+2)){
      for(int i=0; i<m; ++i) for(int j=0; j<n; ++j) D[i][j] =</pre>
      A[i][j];
      for(int i=0; i<m; ++i) B[i] = n+i, D[i][n] = -1, D[i][n+1] =
      for(int j=0; j<n; ++j) N[j] = j, D[m][j] = -c[j];
      N[n] = -1; D[m+1][n] = 1;
  void Pivot(int r, int s) {
   real_t inv = 1/D[r][s];
   for(int i=0; i<m+2; ++i){
      for(int j=0; j<n+2; ++j){
        if(i != r && j != s) D[i][j] -= D[r][j] * D[i][s] * inv;
   }
   for(int i=0; i<m+2; ++i) if(i != r) D[i][s] *= -inv;
   for(int j=0; j<n+2; ++j) if(j != s) D[r][j] *= inv;
   D[r][s] = inv; swap(B[r], N[s]);
  bool Phase(bool p) {
   int x = m + p;
    while(true) {
      int s = -1;
      for(int j=0; j<=n; ++j){
       if(!p && N[j] == -1) continue;
       if(s == -1 || D[x][j] < D[x][s]) s = j;
      if(D[x][s] > -EPS) return true;
      int r = -1;
      for(int i=0; i<m; ++i){</pre>
        if(D[i][s] <= EPS) continue;</pre>
       if(r == -1 || D[i][n+1] / D[i][s] < D[r][n+1] / D[r][s]) r =
       i;
      if(r == -1) return false:
```

```
Pivot(r. s):
 }
 real_t solve(VD &x) {
    int r = 0;
    for(int i=1; i<m; ++i) if(D[i][n+1] < D[r][n+1]) r=i;
    if(D[r][n+1] < -EPS) {
      Pivot(r, n);
      if(!Phase(1) || D[m+1][n+1] < -EPS) return -1/0.0;
      for(int i=0; i<m; ++i) if(B[i] == -1) {
        int s = min_element(D[i].begin(), D[i].end() - 1) -
       D[i].begin();
        Pivot(i, s);
     }
    if(!Phase(0)) return 1/0.0;
    x = VD(n):
    for(int i=0; i<m; ++i) if(B[i] < n) x[B[i]] = D[i][n+1];</pre>
    return D[m][n+1];
 }
};
5.5 Range Prime Counting
// credit :
https://github.com/stjepang/snippets/blob/master/count_primes.cpp
// Primes up to 10^12 can be counted in ~1 second.
const int MAXN = 1000005; // MAXN is the maximum value of sqrt(N) +
bool prime[MAXN];
int prec[MAXN];
vector<int> P;
void init() {
  prime[2] = true;
  for (int i = 3; i < MAXN; i += 2) prime[i] = true;</pre>
```

for (int j = i\*i; j < MAXN; j += i+i) prime[j] = false;</pre>

for (int i = 3; i\*i < MAXN; i += 2){

if (prime[i]){

```
for(int i=1; i<MAXN; i++){</pre>
    if (prime[i]) P.push_back(i);
    prec[i] = prec[i-1] + prime[i];
}
lint rec(lint N, int K) {
  if (N <= 1 | | K < 0) return 0;
  if (N \leq P[K]) return N-1;
  if (N < MAXN && 111 * P[K]*P[K] > N) return N-1 - prec[N] +
  prec[P[K]];
  const int LIM = 250;
  static int memo[LIM*LIM][LIM];
  bool ok = N < LIM*LIM:
  if (ok && memo[N][K]) return memo[N][K];
 lint ret = N/P[K] - rec(N/P[K], K-1) + rec(N, K-1);
  if (ok) memo[N][K] = ret;
  return ret;
}
lint count_primes(lint N) { //less than or equal to
  if (N < MAXN) return prec[N];</pre>
 int K = prec[(int)sqrt(N) + 1];
 return N-1 - rec(N, K) + prec[P[K]];
}
     Discrete Square Root
// https://github.com/tzupengwang/PECaveros/
// blob/master/codebook/math/DiscreteSqrt.cpp
void calcH(int &t, int &h, const int p) {
  int tmp=p-1; for(t=0;(tmp&1)==0;tmp/=2) t++; h=tmp;
// solve equation x^2 \mod p = a
bool solve(int a, int p, int &x, int &y) {
  if (p == 2) { x = y = 1; return true; }
  int p2 = p / 2, tmp = mypow(a, p2, p);
 if (tmp == p - 1) return false;
  if ((p + 1) \% 4 == 0) {
```

```
x=mypow(a,(p+1)/4,p); y=p-x; return true;
 } else {
    int t, h, b, pb; calcH(t, h, p);
    if (t >= 2) {
      do \{b = rand() \% (p - 2) + 2;
     } while (mypow(b, p / 2, p) != p - 1);
      pb = mypow(b, h, p);
   } int s = mypow(a, h / 2, p);
    for (int step = 2; step <= t; step++) {</pre>
     int ss = (((lint)(s * s) % p) * a) % p;
     for(int i=0;i<t-step;i++) ss=(lint)ss*ss%p;;</pre>
     if (ss + 1 == p) s = (s * pb) % p;
      pb = ((lint)pb * pb) % p;
   x = ((lint)s * a) % p; y = p - x;
 } return true:
     Miller-Rabin Test + Pollard Rho Factorization
namespace miller_rabin{
 lint mul(lint a, lint b, lint p){
   lint ret = 0;
    while(a){
     if(a\&1) ret = (ret + b) \% p;
     a >>= 1;
      b = (b << 1) \% p;
    return ret;
```

lint ipow(lint x, lint y, lint p){

if(y&1) ret = mul(ret, piv, p);

lint ret = 1, piv = x % p;

piv = mul(piv, piv, p);

bool miller\_rabin(lint x, lint a){

if(x % a == 0) return 0;

}

while(y){

v >>= 1;

return ret;

```
lint d = x - 1;
    while(1){
      lint tmp = ipow(a, d, x);
      if(d&1) return (tmp != 1 && tmp != x-1);
      else if(tmp == x-1) return 0;
      d >>= 1;
   }
  bool isprime(lint x){
   for(auto &i : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}){
      if(x == i) return 1;
     if (x > 40 \&\& miller_rabin(x, i)) return 0;
    if(x \le 40) return 0;
    return 1;
}
namespace pollard_rho{
  lint f(lint x, lint n, lint c){
    return (c + miller_rabin::mul(x, x, n)) % n;
  void rec(lint n, vector<lint> &v){
   if(n == 1) return;
   if(n \% 2 == 0){
   v.push_back(2);
    rec(n/2, v);
    return;
  if(miller_rabin::isprime(n)){
   v.push_back(n);
    return;
  lint a, b, c;
  while(1){
   a = rand() \% (n-2) + 2;
   b = a;
    c = rand() \% 20 + 1;
    do{
```

```
a = f(a, n, c);
b = f(f(b, n, c), n, c);
}while(gcd(abs(a-b), n) == 1);
if(a != b) break;
}
lint x = gcd(abs(a-b), n);
rec(x, v);
rec(n/x, v);
}
vector<lint> factorize(lint n){
  vector<lint> ret;
  rec(n, ret);
  sort(ret.begin(), ret.end());
  return ret;
}
};
```

# 5.8 Highly Composite Numbers, Large Prime

< 10	^k number	divisors	2 3 5 71113171923293137
1	6	4	1 1
2	60	12	2 1 1
3	840	32	3 1 1 1
4	7560	64	3 3 1 1
5	83160	128	3 3 1 1 1
6	720720	240	4 2 1 1 1 1
7	8648640	448	6 3 1 1 1 1
8	73513440	768	5 3 1 1 1 1 1
9	735134400	1344	6 3 2 1 1 1 1
10	6983776800	2304	5 3 2 1 1 1 1 1
11	97772875200	4032	6 3 2 2 1 1 1 1
12	963761198400	6720	6 4 2 1 1 1 1 1 1
13	9316358251200	10752	6 3 2 1 1 1 1 1 1 1
14	97821761637600	17280	5 4 2 2 1 1 1 1 1 1
15	866421317361600	26880	6 4 2 1 1 1 1 1 1 1 1
16	8086598962041600	41472	8 3 2 2 1 1 1 1 1 1 1
17	74801040398884800	64512	6 3 2 2 1 1 1 1 1 1 1 1
18	897612484786617600	103680	8 4 2 2 1 1 1 1 1 1 1 1

< 10^1	x prime	# of prime	< 10	k prime
1	7	4	10	999999967
2	97	25	11	9999999977
3	997	168	12	999999999989
4	9973	1229	13	999999999971
5	99991	9592	14	9999999999973
6	999983	78498	15	99999999999989
7	9999991	664579	16	99999999999937
8	99999989	5761455	17	999999999999999
9	99999937	50847534	18	9999999999999999

#### NTT Prime:

```
998244353 = 119 \times 2^{23} + 1. Primitive root: 3. 985661441 = 235 \times 2^{22} + 1. Primitive root: 3. 1012924417 = 483 \times 2^{21} + 1. Primitive root: 5.
```

# 6 Miscellaneous

# 6.1 Popular Optimization Technique

- Convex Hull Trick (cf: 5.4 for dynamic slopes)
- Divide and Conquer Optimization
- Knuth's  $O(n^2)$  Optimal BST: minimize  $D_{i,j} = Min_{i \leq k < j}(D_{i,k} + D_{k+1,j}) + C_{i,j}$ . Quadrangle Inequality:  $C_{a,c} + C_{b,d} \leq C_{a,d} + C_{b,c}$ ,  $C_{b,c} \leq C_{a,d}$ . Now monotonicity holds.
- Sqrt batch processing Save queries in buffer, and update in every sqrt steps (cf : IOI 2011 Elephant. hyea calls it "ainta technique")
- Dynamic insertion in static set (Make O(lgn) copy. Merge like binomial heap.)
- Offline insertion / deletion in insert-only set (Pair insertion-deletion operation, and regard it as range query)
- Mo's algorithm trick (on tree)
- Aliens trick: Partition n elements into k contiguous interval: Partition naively without k restriction. Penalize / Reward the partition by changing cost function. Mix this with binary search.

# 6.2 Bit Twiddling Hack

return x;

```
int __builtin_clz(int x);//number of leading zero
int __builtin_ctz(int x);//number of trailing zero
int __builtin_clzll(long long x);//number of leading zero
int __builtin_ctzll(long long x);//number of trailing zero
int __builtin_popcount(int x);// number of 1-bits in x
int __builtin_popcountll(long long x);//number of 1-bits in x
lsb(n): (n & -n); // last bit (smallest)
floor(log2(n)): 31 - __builtin_clz(n | 1);
floor(log2(n)): 63 - __builtin_clzll(n | 1);
//compute next perm. ex) 00111, 01011, 01101, 01110, 10011, 10101...
long long next_perm(long long v){
 long long t = v \mid (v-1);
 return (t + 1) \mid (((^t \& -^t) - 1) >> (builtin ctz(v) + 1));
     Fast Integer IO
static char buf[1 << 19]; // size : any number geq than 1024
static int idx = 0;
static int bytes = 0;
static inline int _read() {
    if (!bytes || idx == bytes) {
        bytes = (int)fread(buf, sizeof(buf[0]), sizeof(buf), stdin);
        idx = 0:
   }
    return buf[idx++]:
static inline int readInt() {
    int x = 0, s = 1;
    int c = _read();
    while (c \le 32) c = read();
    if (c == '-') s = -1, c = _read();
    while (c > 32) x = 10 * x + (c - '0'), c = _read();
    if (s < 0) x = -x;
```

# 6.4 OSRank in g++

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
typedef
tree<int, null_type, less<int>, rb_tree_tag,
tree_order_statistics_node_update> ordered_set;
ordered set X:
X.insert(1); X.insert(2); X.insert(4); X.insert(8); X.insert(16);
cout<<*X.find_by_order(1)<<endl; // 2</pre>
cout<<*X.find_by_order(2)<<endl; // 4</pre>
cout<<*X.find_by_order(4)<<endl; // 16</pre>
cout<<(end(X)==X.find_by_order(6))<<endl; // true</pre>
cout<<X.order_of_key(-5)<<endl; // 0</pre>
cout<<X.order_of_key(1)<<endl; // 0</pre>
cout<<X.order_of_key(3)<<endl; // 2</pre>
cout<<X.order_of_key(4)<<endl; // 2</pre>
cout<<X.order_of_key(400)<<endl; // 5</pre>
```

# 6.5 Nasty Stack Hacks

# 6.6 C++ / Environment Overview

```
// vimrc : set nu sc ci si ai sw=4 ts=4 bs=2 mouse=a syntax on
// compile : g++ -o PROB PROB.cpp -std=c++11 -Wall -02
// options : -fsanitize=address -Wfatal-errors
#include <bits/stdc++.h> // magic header
using namespace std; // magic namespace
// how to use rand (in 2017)
mt19937 rng(0x14004);
int randint(int lb, int ub){ return
uniform_int_distribution<int>(lb, ub)(rng); }
// comparator overload
auto cmp = [](seg a, seg b){return a.func() < b.func(); };</pre>
set<seg, decltype(cmp)> s(cmp);
map<seg, int, decltype(cmp)> mp(cmp);
priority_queue<seg, vector<seg>, decltype(cmp)> pq(cmp); // max heap
// hash func overload
struct point{
int x, y;
bool operator==(const point &p)const{ return x == p.x && y == p.y; }
struct hasher {
size_t operator()(const point &p)const{ return p.x * 2 + p.y * 3; }
unordered_map<point, int, hasher> hsh;
```