Breaking the Nonsmooth Barrier: A Scalable Parallel Method for Composite Optimization

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Summary

Optimization methods need to be adapted to the parallel setting to leverage modern computer architectures.

Highly efficient variants of stochastic gradient descent have been recently proposed, such as Hogwild [1], Kromagnon [2], ASAGA [3].

They assume that the objective function is smooth, so are inapplicable to problems such as Lasso, optimization with constraints, etc.

Contributions:

- . Sparse Proximal SAGA, a sparse variant of the linearly-convergent proximal SAGA algorithm.
- . ProxASAGA, the first parallel asynchronous variance-reduced method that supports composite objective functions.

Problem Setting

Objective: develop parallel asynchronous method for problems of the form $\min_{oldsymbol{x} \in \mathbb{R}^p} \lim_{i \in \mathbb{R}^p} \frac{1}{n} \sum_{i=1}^n f_i(oldsymbol{x}) + h(oldsymbol{x}) ,$

- f_i is differentiable with L-Lipschitz gradient.
- h is block-separable $(h(x) = \sum_B h_B([{m x}]_B))$ and "simple" in the sense that we have access to $\mathbf{prox}_{\gamma h} \stackrel{\mathsf{def}}{=} \mathbf{arg} \, \mathbf{min}_{m{x}} \, \gamma h(m{x}) + \frac{1}{2} \|m{x} - m{z}\|^2$.
- includes Lasso, group Lasso or ERM with box constraints.

Variance-reduced stochastic gradient methods are natural candidates due to their state of the art performance and recent asynchronous variants.

The **SAGA** algorithm [4] maintains current iterate $oldsymbol{x} \in \mathbb{R}^p$ and table of historical gradients $oldsymbol{lpha} \in \mathbb{R}^{n imes p}$. At each iteration, it samples an index $i \in$ $\{1,\ldots,n\}$ and computes next iterate $({m x}^+,{m lpha}^+)$ as

$$\boldsymbol{x}^{+} = \mathbf{prox}_{\gamma g}(\boldsymbol{x} - \gamma(\nabla f_i(\boldsymbol{x}) - \boldsymbol{\alpha}_i + \overline{\boldsymbol{\alpha}})), \ \boldsymbol{\alpha}_i^{+} = \nabla f_i(\boldsymbol{x}).$$
 (SAGA)

Difficulty of a Composite Extension

- Existing methods exhibit best performance when updates are sparse.
- Even in the presence of sparse gradients, the (SAGA) update is not sparse due to the presence of $\overline{\alpha}$ and prox.
- Existing convergence proofs bound noise from asynchrony using by the Lipschitz of the gradient. Property does not extend to composite case.

A new sequential algorithm: Sparse Proximal SAGA

The algorithm relies on the following quantities

• Extended support T_i : set of blocks that intersect with ∇f_i .

$$T_i \stackrel{\mathsf{def}}{=} \{B : \mathsf{supp}(\nabla f_i) \cap B \neq \varnothing, \ B \in \mathcal{B}\}$$

- For each block $B \in \mathcal{B}$, $d_B \stackrel{\mathsf{def}}{=} n/n_B$, where $n_B := \sum_i \mathbb{1}\{B \in T_i\}$ is the number of times that $B \in T_i$.
- $m{D}_i$ is a diagonal matrix defined block-wise $[m{D}_i]_{B,B} \stackrel{\mathsf{def}}{=} d_B \mathbb{1}\{B \in T_i\} m{I}_{|B|}$.
- ullet φ_i is a block-wise reweighting of h: $arphi_i \stackrel{ ext{def}}{=} \sum_{B \in T_i} d_B h_B(oldsymbol{x})$

Justification. The following properties are verified

 $\boldsymbol{D}_i \boldsymbol{x}$ is zero outside T_i $\varphi_i(\boldsymbol{x})$ is zero outside T_i (sparsity) $\mathbb{E}_i \, oldsymbol{D}_i = oldsymbol{I}$ (unbiasedness) $\mathbb{E}_i \, \varphi_i = h$

Algorithm. As SAGA, it maintains current iterate $oldsymbol{x} \in \mathbb{R}^p$ and table of historical gradients $\pmb{\alpha} \in \mathbb{R}^{n \times p}$. At each iteration, it samples an index $i \in \{1,\ldots,n\}$ and computes next iterate $(\boldsymbol{x}^+,\boldsymbol{\alpha}^+)$ as

$$oldsymbol{v}_i =
abla f_i(oldsymbol{x}) - oldsymbol{lpha}_i + oldsymbol{D}_i \overline{oldsymbol{lpha}}; \ oldsymbol{x}^+ = \mathbf{prox}_{\gamma arphi_i} \left(oldsymbol{x} - \gamma oldsymbol{v}_i
ight); \ oldsymbol{lpha}_i^+ =
abla f_i(oldsymbol{x})$$

Features

- Iteration cost $\mathcal{O}(|T_i|)$ (only coefficients in T_i are updated)
- Easy to implement (compared to the lagged update approach [5]).
- Parallelizable.

Convergence Analysis

For step size $\gamma = \frac{1}{5L}$ and f μ -strongly convex ($\mu > 0$), Sparse Proximal SAGA converges geometrically in expectation. At iteration t we have

$$\mathbb{E}\|\boldsymbol{x}_t - \boldsymbol{x}^*\|^2 \le (1 - \frac{1}{5}\min\{\frac{1}{n}, \frac{1}{\kappa}\})^t C_0,$$

(SAGA) with $C_0 = \| \boldsymbol{x}_0 - \boldsymbol{x}^* \|^2 + \frac{1}{5L^2} \sum_{i=1}^n \| \boldsymbol{\alpha}_i^0 - \nabla f_i(\boldsymbol{x}^*) \|^2$ and $\kappa = \frac{L}{n}$ (condition

Implications

- Adaptivity to strong convexity, i.e., no need to know strong convexity parameter to obtain linear convergence.
- In the "big data regime" $(n \ge \kappa)$: convergence rate is $\mathcal{O}(1/n)$.
- In the "ill-conditioned regime" ($n \le \kappa$): convergence rate is $\mathcal{O}(1/\kappa)$.

A New Parallel Algorithm: Proximal Asynchronous SAGA

The Proximal Asynchronous SAGA (ProxASAGA) runs Inconsistent read: no locks while reading, i.e., read the vector while another core might be writing to it.

Algorithm 1 ProxASAGA

```
Initialize shared variables x, (\alpha_i)_{i=1}^n, \overline{\alpha}
keep doing in parallel
          Sample i uniformly in \{1, ..., n\}
           [\hat{\boldsymbol{x}}]_{T_i} = \text{inconsistent read of } \boldsymbol{x} \text{ on } T_i
           \hat{m{lpha}}_i = 	ext{inconsistent read of } m{lpha}_i
            [\overline{\alpha}]_{T_i} = \text{inconsistent read of } \overline{\alpha} \text{ on } T_i
            [\delta oldsymbol{lpha}]_{S_i} = [
abla f_i(\hat{oldsymbol{x}})]_{S_i} - [\hat{oldsymbol{lpha}}_i]_{S_i}
           [\hat{oldsymbol{v}}]_{T_i} = [\delta oldsymbol{lpha}]_{T_i} + [oldsymbol{D}_i \overline{oldsymbol{lpha}}]_{T_i}
           [\delta oldsymbol{x}]_{T_i} = [\mathbf{prox}_{\gamma_{\mathcal{O}_i}}(\hat{oldsymbol{x}} - \gamma \hat{oldsymbol{v}})]_{T_i} - [\hat{oldsymbol{x}}]_{T_i}
          for B in T_i do
                   for b in B do
                            [\boldsymbol{x}]_b \leftarrow [\boldsymbol{x}]_b + [\delta \boldsymbol{x}]_b
                                                                                                                                              > atomic
                            if b \in \text{support}(\nabla f_i) then
                                     [\overline{\boldsymbol{\alpha}}]_b \leftarrow [\overline{\boldsymbol{\alpha}}]_b + 1/n[\delta \boldsymbol{\alpha}]_b
                                                                                                                                              end if
                   end for
          end for
```

Perturbed iterates framework

19: end parallel loop

Problem: Analysis of parallel algorithms is hard.

 $\alpha_i \leftarrow \nabla f_i(\hat{\boldsymbol{x}})$ (scalar update)

Solution: Cast them as sequential algorithms working on *perturbed* inputs. Distinguish:

- $\hat{m{x}}_t$: inconsistent quantity read by the cores
- x_t : the *virtual* iterate defined by

 $m{x}_{t+1} \stackrel{\mathsf{def}}{=} m{x}_t - \gamma m{g}_t$ with $m{g}(m{x}, m{v}, i) = rac{1}{\gamma} (\hat{m{x}}_t - \mathbf{prox}_{\gamma arphi_i} (\hat{m{x}}_t - \gamma \hat{m{v}}_{i_t}))$, Interpret \hat{x}_t as a noisy version of \boldsymbol{x}_t due to asynchrony.

Analysis preliminaries

Definition (measure of sparsity). Let $\Delta := \max_{B \in \mathcal{B}} |\{i : T_i \ni B\}|/n$. This is the normalized maximum number of times that a block appears in the extended support. We always have $1/n \le \Delta \le 1$.

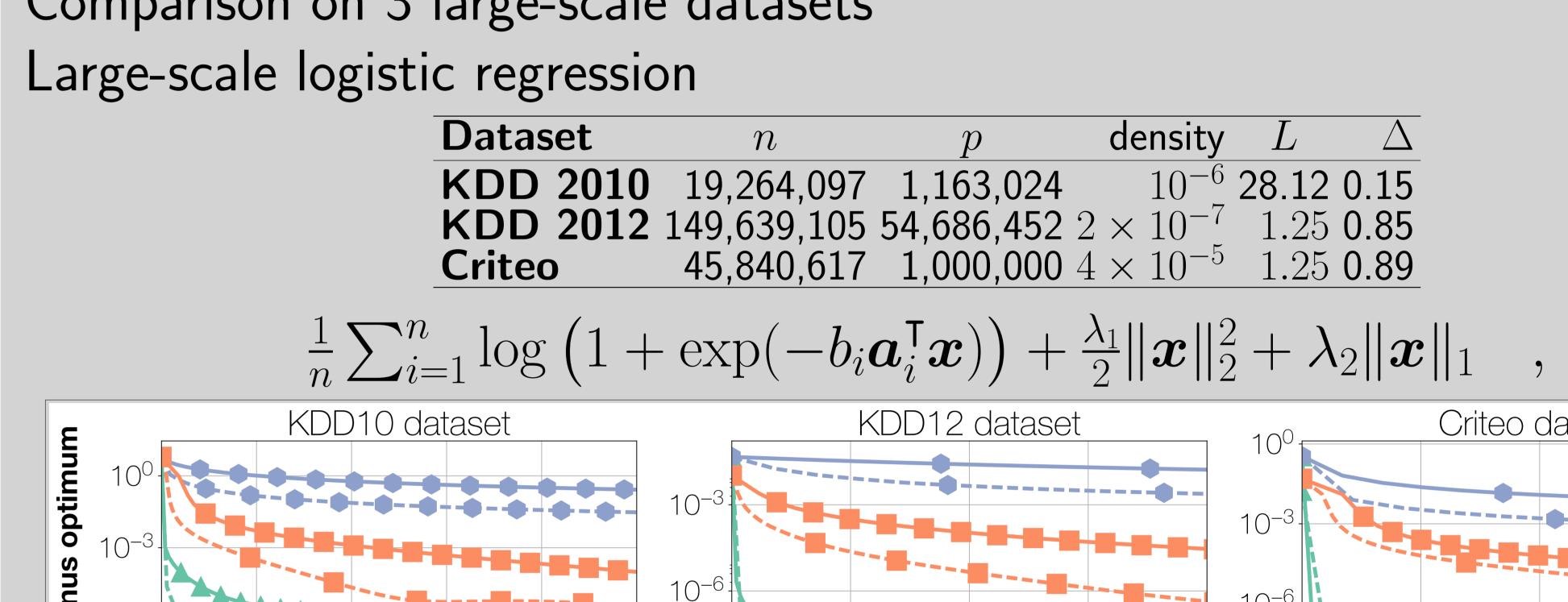
Definition (delay bound). τ is a uniform bound on the maximum delay between two iterations processed concurrently.

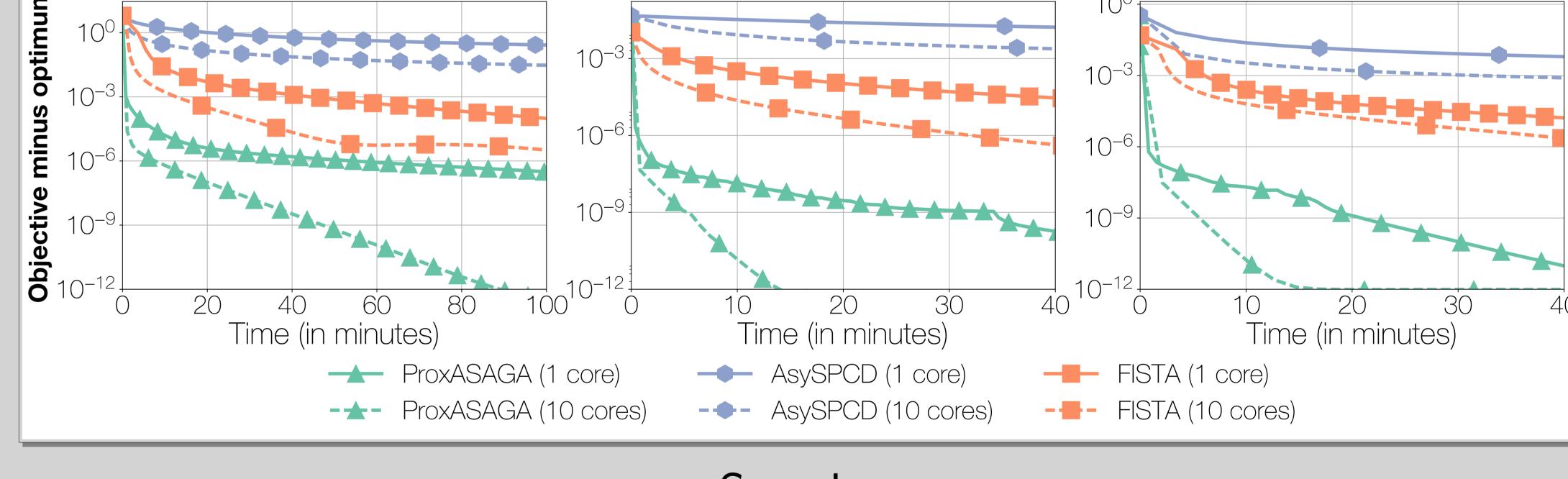
Convergence guarantee of ProxASAGA

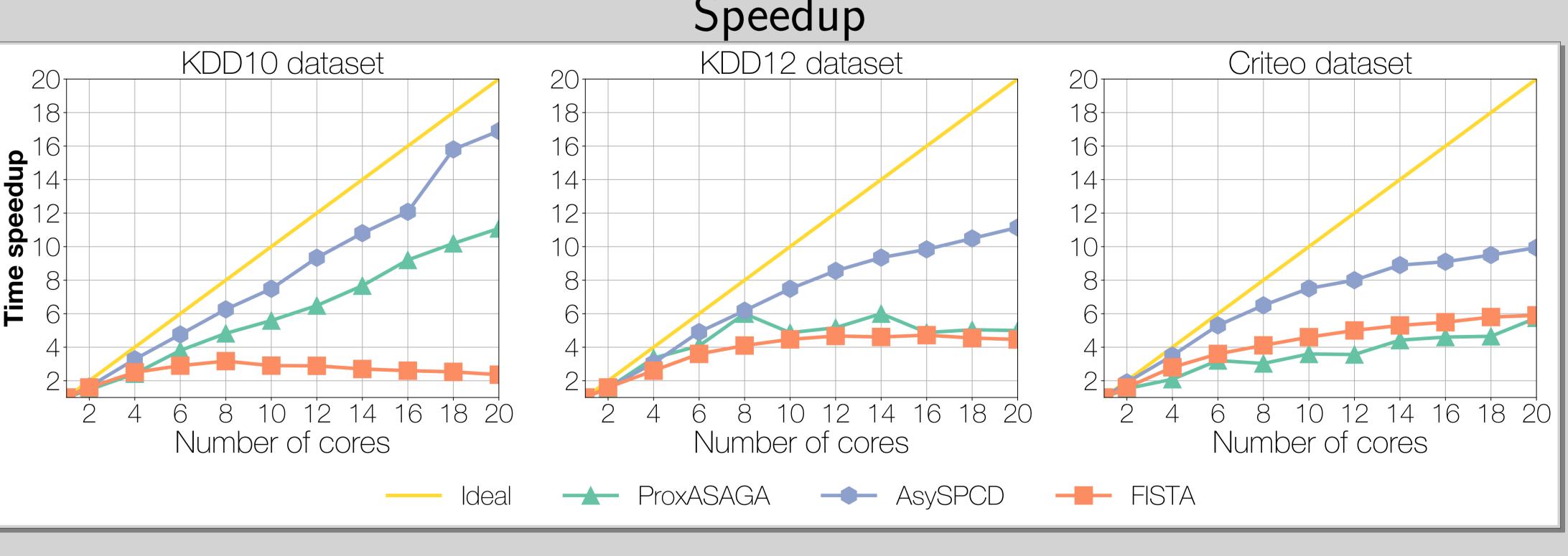
Suppose $\tau \leq \frac{1}{10\sqrt{\Lambda}}$. For any step size $\gamma = \frac{a}{L}$ with $a \leq a^*(\tau) := \frac{1}{36} \min\{1, \frac{6\kappa}{\tau}\}$, the inconsistent read iterates of ProxASAGA converge in expectation at a geometric rate factor of at least: $\rho(a) = \frac{1}{5}\min\left\{\frac{1}{n},a_{\kappa}^{1}\right\}$, i.e. $\mathbb{E}\|\hat{x}_{t}-x^{*}\|^{2} \leq 1$ $(1ho)^t \tilde{C}_0$, where \tilde{C}_0 is a constant independent of t ($pprox rac{n\kappa}{a} C_0$ with C_0 as defined in Theorem ??).

Experimental results

Comparison on 3 large-scale datasets







References

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