

1.1 Pearson Correlator 1D Representation

OMG a 2D correlator equation is huge

To translate the formula into code we first need to understand it.

Input: template $h[n] \sim 3: [2, 3, 5]$
discrete input $x[n] \sim 10: [1, 2, 3, 5, 4, 4, 2, 0, 1, 12, 13, 15]$

Output: sequence of pearson correlations $y[n] \sim 8: [.....]$

Formula:
$$y[n] = \frac{\sum_{k=0}^{L-1} (x[n+k] - \bar{x})(h[k] - \bar{h})}{\sqrt{\sum_{k=0}^{L-1} (x[n+k] - \bar{x})^2} \cdot \sqrt{\sum_{k=0}^{L-1} (h[k] - \bar{h})^2}}$$

Simplified:
$$y[n] = \frac{L(\sum_k x[n+k]h[k]) - (\sum_k x[n+k])(\sum_k h[k])}{\sqrt{L\sum_k x[n+k]^2 - (\sum_k x[n+k])^2} \cdot \sqrt{L\sum_k h[k]^2 - (\sum_k h[k])^2}}$$

2: [1, 42]

10: [1, 42, 1, 42, 1, 42, 10, 52, 10, 52]

$$\text{sum} = 42 \cdot 1 + 42 \cdot 42$$

$$\text{sum1} = 42 + 42$$

$$\text{sum2} = 1 + 42$$

this is a lil note about some code for my assignment.

1.2 Discrete Fourier Transform w/ Vandermonde matrix

$$W = \begin{bmatrix} 1 & 1 & 1 & 1 & \dots & 1 \\ 1 & \omega & \omega^2 & \omega^3 & \dots & \omega^{N-1} \\ 1 & \omega^2 & \omega^4 & \omega^6 & \dots & \omega^{2(N-1)} \\ 1 & \omega^3 & \omega^6 & \omega^9 & \dots & \omega^{3(N-1)} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \omega^{N-1} & \omega^{2(N-1)} & \omega^{3(N-1)} & \dots & \omega^{(N-1)(N-1)} \end{bmatrix}$$

DFT converts $x[n]$ from time to frequency domain

$$X[k] = \sum_{n=0}^{N-1} x[n] \cdot e^{-j \frac{2\pi}{N} kn}$$

k : Frequency index

N : Samples in $x[n]$:: $\text{len}(x)$

$X[k]$: complex rep of Amp and Phase

$$\omega = e^{-j \frac{2\pi}{N}} = \cos\left(\frac{2\pi}{N}\right) + j \sin\left(\frac{2\pi}{N}\right)$$

(a, b) i mean this a lil nugget that you can use to make something similar to math PS. I use this for 2.3

1.3 Inverse Discrete Fourier Transform w/ Vandermonde matrix

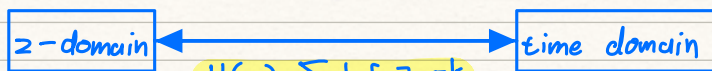
$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi}{N} kn}$$

I am now unsure if this is actually difficult or simply just a different Omega?

So, my guess was correct, but its not just a different Omega. The second difference is that each $x[n]$ element is divided by N

$$\omega = e^{j \frac{2\pi}{N} \cdot kn} = \cos\left(\frac{2\pi}{N}\right) + j \sin\left(\frac{2\pi}{N}\right)$$

Z-Domain to time-Domain



$$H(z) = \sum_k h[n] z^{-k}$$

this was apparently not it!!

Useless: (

FIR Filter ~ Roots/zeros of system
 $X(z)$

→ FIR Filter Representation in three domains

1) TIME: $h[n] = \sum_k b_k \delta[n-k]$

2) FREQUENCY: $H(e^{j\omega}) = \sum_k h[n] e^{-j\omega k}$
 or $H(e^{j\omega}) = H(z) \big|_{z=e^{j\omega}}$

3) Z-DOMAIN: $H(z) = \sum_k h[n] z^{-k}$

time ↔ z-domain via $H(e^{j\omega})$

Input: M : number of roots

$\theta_k \times M$: Roots $\{z_k\}$ as angles θ_k ; $z_k = e^{j\theta_k}$

Output: -time-domain impulse response $h[n]$

constant factor = 1

-length of $h[n]$ should be $M+1$

z-transform of a filter:

$$H(z) = G \prod_{k=1}^M (1 - z_k \cdot z^{-1}) = G \prod_{k=1}^M \left(\frac{z - z_k}{z} \right)$$

z_k = Zeros of the system
 $G = 1$

Examples: 1

1.047198; -1.047198

→ $(1 - (0.5 + 0.866j)z^{-1})(1 - (-0.5 + 0.866j)z^{-1})$

$1 + (0.5 + 0.866j)(-0.5 + 0.866j)$

$1 + (-0.25 + 0.5 \times 0.8j - 0.5 \times 0.8j - 0.866j \times 0.866j) z^{-2}$

$1 - 1z^{-2}$

$\alpha = 0.5 + 0.866j$

1) convert angles (θ_k) into roots (z_k).

$z_k = e^{j\theta_k} = \cos(\theta_k) + j \cdot \sin(\theta_k)$

for simplicity we will keep calling these roots z_k

2) Sub into $H(z) = G \prod_{k=1}^M (1 - z_k \cdot z^{-1})$ and expand. # this is the difficult part.
 # G in this assignment is 1

Let's use an example where we have 2 zeros

$H(z) = (1 - z_1 \cdot z^{-1})(1 - z_2 \cdot z^{-1})$
 $= 1 - z_2 z^{-1} - z_1 z^{-1} + z_1 \cdot z_2 \cdot z^{-2}$
 $= 1 - (z_1 + z_2) z^{-1} + (z_1 \cdot z_2) \cdot z^{-2}$

$\therefore h[n] = [1; -(z_1 + z_2); (z_1 \cdot z_2)]$
 if $M = 2$

Ok... 3 zeros:

$H(z) = (1 - z_1 \cdot z^{-1})(1 - z_2 \cdot z^{-1})(1 - z_3 \cdot z^{-1})$
 $= [1 - (z_1 + z_2) z^{-1} + (z_1 \cdot z_2) \cdot z^{-2}] [1 - z_3 \cdot z^{-1}]$
 $= 1 - z_3 z^{-1} - (z_1 + z_2) z^{-1} + (z_1 + z_2) z^{-1} \cdot z_3 z^{-1} + z_1 \cdot z_2 \cdot z^{-2} - z_1 \cdot z_2 \cdot z^{-2} \cdot z_3 z^{-1}$
 $= 1 - z_3 z^{-1} - z_1 z^{-1} - z_2 z^{-1} + z_1 \cdot z_3 \cdot z^{-2} + z_2 \cdot z_3 \cdot z^{-2} + z_1 \cdot z_2 \cdot z^{-2} - z_1 \cdot z_2 \cdot z_3 \cdot z^{-3}$
 $= 1 - (z_1 + z_2 + z_3) z^{-1} + (z_1 \cdot z_3 + z_2 \cdot z_3 + z_1 \cdot z_2) \cdot z^{-2} - (z_1 \cdot z_2 \cdot z_3) z^{-3}$

$\therefore h[n] = [1; -(z_1 + z_2 + z_3); (z_1 \cdot z_3 + z_2 \cdot z_3 + z_1 \cdot z_2); -(z_1 \cdot z_2 \cdot z_3)]$

→ To translate this into code, I had the following idea of using an array for coefficients.

↳ Start with array of roots $[\cos(\theta_k) + j \sin(\theta_k); \dots]$

↳ have an array of coefficients... coeffs = []

↳ append the first "1" element

The idea of the coeff array is to have the indexes represent powers:
 1: $z^0 = 1$ 2: z^{-1} 3: z^{-2}

I figured it out Bitch!!!
weep weep!

Lets use 10 roots:

$$\begin{aligned} a_{10} &= 1 && \text{(assuming monic polynomial)} \\ a_9 &= -(r_1 + r_2 + r_3 + \dots + r_{10}) && \text{(sum of single roots)} \\ a_8 &= \sum_{i < j} r_i r_j && \text{(sum of products of pairs of roots)} \\ a_7 &= (-1) \sum_{i < j < k} r_i \cdot r_j \cdot r_k && \text{(sum of all triplets of roots)} \\ a_6 &= && \text{the pattern should be clear here.} \end{aligned}$$

Now for code we need 3 functions:

- 1) main function okay
- 2) compute combinations difficult NAH `itertools.combinations(array, n-root)`
- 3) product of combinations. easy

→ Well, this would have been great untill we came to a time complexity of $O(n^3)$.

↪ I found a new method!

Horner's Method → this algorithm is based on horner's rule in which a polynomial is written in nested form.