

Massachusetts Institute of Technology

MIT NULL

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Contest (1)

template.cpp

55 line

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
typedef long double ld;
typedef double db;
typedef string str;
typedef pair<int, int> pi;
typedef pair<11,11> pl;
typedef pair<ld,ld> pd;
typedef complex<ld> cd;
typedef vector<int> vi;
typedef vector<ll> vl;
typedef vector<ld> vd;
typedef vector<str> vs;
typedef vector<pi> vpi;
typedef vector<pl> vpl;
typedef vector<cd> vcd;
#define FOR(i,a,b) for (int i = (a); i < (b); ++i)
#define FOR(i,a) FOR(i,0,a)
#define ROF(i,a,b) for (int i = (b)-1; i \ge (a); --i)
#define ROF(i,a) ROF(i,0,a)
#define trav(a,x) for (auto& a : x)
#define mp make_pair
#define pb push back
#define eb emplace_back
#define f first
#define s second
#define lb lower bound
#define ub upper bound
#define sz(x) (int)x.size()
\#define all(x) begin(x), end(x)
#define rall(x) rbegin(x), rend(x)
#define rsz resize
#define ins insert
const int MOD = 1e9+7; // 998244353 = (119 << 23) +1
const 11 INF = 1e18;
const int MX = 2e5+5;
const ld PI = 4*atan((ld)1);
template<class T> bool ckmin(T& a, const T& b) { return a > b ?
  \hookrightarrow a = b, 1 : 0; }
template<class T> bool ckmax(T& a, const T& b) { return a < b ?</pre>
  \hookrightarrow a = b, 1 : 0; }
mt19937 rng(chrono::steady_clock::now().time_since_epoch().

→count());
int main() {
    cin.sync_with_stdio(0); cin.tie(0);
.bashrc
    g++ -std=c++11 -O2 -Wall -Wl,-stack_size -Wl,0x10000000 -o
       $1 $1.cc
```

```
run() {
    co $1 && ./$1
.vimrc
set cin aw ai is ts=4 sw=4 tm=50 nu noeb ru cul
sy on | im jk <esc> | im kj <esc>
set mouse=a
set ww+=<,>,[,]
hash.sh
# Hashes a file, ignoring all whitespace and comments. Use for
# verifying that code was correctly typed.
cpp -dD -P -fpreprocessed | tr -d '[:space:]' | md5sum |cut -c-6
troubleshoot.txt
Pre-submit:
Write a few simple test cases, if sample is not enough.
Are time limits close? If so, generate max cases.
Is the memory usage fine?
Could anything overflow?
Make sure to submit the right file.
Wrong answer:
Print your solution! Print debug output, as well.
Are you clearing all datastructures between test cases?
Can your algorithm handle the whole range of input?
Read the full problem statement again.
Do you handle all corner cases correctly?
Have you understood the problem correctly?
Any uninitialized variables?
Anv overflows?
Confusing N and M, i and j, etc.?
Are you sure your algorithm works?
What special cases have you not thought of?
Are you sure the STL functions you use work as you think?
Add some assertions, maybe resubmit.
Create some testcases to run your algorithm on.
Go through the algorithm for a simple case.
Go through this list again.
Explain your algorithm to a team mate.
Ask the team mate to look at your code.
Go for a small walk, e.g. to the toilet.
Is your output format correct? (including whitespace)
Rewrite your solution from the start or let a team mate do it.
Runtime error:
Have you tested all corner cases locally?
Any uninitialized variables?
Are you reading or writing outside the range of any vector?
Any assertions that might fail?
Any possible division by 0? (mod 0 for example)
Any possible infinite recursion?
Invalidated pointers or iterators?
Are you using too much memory?
Debug with resubmits (e.g. remapped signals, see Various).
Time limit exceeded:
Do you have any possible infinite loops?
What is the complexity of your algorithm?
Are you copying a lot of unnecessary data? (References)
How big is the input and output? (consider scanf)
Avoid vector, map. (use arrays/unordered_map)
What do your team mates think about your algorithm?
Memory limit exceeded:
```

What is the max amount of memory your algorithm should need? Are you clearing all datastructures between test cases?

Data Structures (2)

2.1 STL

MapComparator.h

Description: custom comparator for map / set

d0cc31, 8 lines

```
struct cmp {
  bool operator()(const int& 1, const int& r) const {
    return 1 > r;
  }
};

set<int,cmp> s; // FOR(i,10) s.insert(rand()); trav(i,s) ps(i);
map<int,int,cmp> m;
```

CustomHash.h

Description: faster than standard unordered map

e7c12c, 23 lines

```
struct chash {
 static uint64 t splitmix64(uint64 t x) {
    // http://xorshift.di.unimi.it/splitmix64.c
   x += 0x9e3779b97f4a7c15;
   x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
   x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
   return x ^ (x >> 31);
 size_t operator()(uint64_t x) const {
   static const uint64_t FIXED_RANDOM =
     chrono::steady_clock::now()
      .time_since_epoch().count();
    return splitmix64(x + FIXED_RANDOM);
};
template<class K, class V> using um = unordered_map<K, V, chash</pre>
template < class K, class V> using ht = qp_hash_table < K, V, chash
  ⇒>:
template<class K, class V> V get(ht<K,V>& u, K x) {
 return u.find(x) == end(u) ? 0 : u[x];
```

OrderStatisticTree.h

Description: A set (not multiset!) with support for finding the n'th element, and finding the index of an element. To get a map, change null-type. **Time:** $\mathcal{O}(\log N)$

```
<ext/pb.ds/tree.policy.hpp>, <ext/pb.ds/assoc.container.hpp> c5d6f2, 17 lines
using namespace __gnu_pbds;

template <class T> using Tree = tree<T, null_type, less<T>,
    rb_tree_tag, tree_order_statistics_node_update>;

#define ook order_of_key
#define fbo find_by_order

void treeExample() {
    Tree<int> t, t2; t.insert(8);
    auto it = t.insert(10).f;
    assert(it == t.lb(9));
    assert(t.ook(10) == 1);
```

```
assert(t.ook(11) == 2);
assert(*t.fbo(0) == 8);
t.join(t2); // assuming T < T2 or T > T2, merge t2 into t
}
```

Rope.h

Description: insert element at n-th position, cut a substring and re-insert somewhere else

Time: $\mathcal{O}(\log N)$ per operation?

LineContainer.h

Description: Given set of lines, computes greatest y-coordinate for any x **Time:** $\mathcal{O}(\log N)$

```
mutable 11 k, m, p; // slope, y-intercept, last optimal x
  11 eval (11 x) { return k*x+m; }
  bool operator<(const Line& o) const { return k < o.k; }</pre>
  bool operator<(11 x) const { return p < x; }
struct LC : multiset<Line,less<>>> {
  // for doubles, use inf = 1/.0, div(a,b) = a/b
  const ll inf = LLONG MAX:
  ll div(ll a, ll b) { return a/b-((a^b) < 0 \&\& a\%b); } //
     \hookrightarrowfloored division
  ll bet(const Line& x, const Line& y) { // last x such that
     \hookrightarrow first line is better
    if (x.k == y.k) return x.m >= y.m? inf : -inf;
    return div(y.m-x.m,x.k-y.k);
  bool isect(iterator x, iterator y) { // updates x->p,
     \hookrightarrowdetermines if v is unneeded
    if (y == end()) { x->p = inf; return 0; }
    x->p = bet(*x,*y); return x->p >= y->p;
  void add(ll k, ll m) {
    auto z = insert(\{k, m, 0\}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() \&\& isect(--x, y)) isect(x, y = erase(y));
    while ((y = x) != begin() \&\& (--x)->p >= y->p) isect(x,
       \rightarrowerase(v));
  ll query(ll x) {
    assert(!empty());
    auto 1 = *lower_bound(x);
    return 1.k * x + 1.m;
};
```

2.2 1D Range Queries

RMQ.h

Description: 1D range minimum query **Time:** $O(N \log N)$ build, O(1) query

0a1f4a, 25 lines

```
template<class T> struct RMQ {
 constexpr static int level(int x) {
   return 31-__builtin_clz(x);
 } // floor(log_2(x))
 vector<vi> jmp;
 vector<T> v;
 int comb(int a, int b) {
   return v[a] == v[b] ? min(a,b) : (v[a] < v[b] ? a : b);
 } // index of minimum
 void init(const vector<T>& _v) {
   v = v; imp = \{vi(sz(v))\}; iota(all(jmp[0]), 0);
   for (int j = 1; 1 << j <= sz(v); ++j) {
     jmp.pb(vi(sz(v)-(1<< j)+1));
     FOR(i, sz(jmp[j])) jmp[j][i] = comb(jmp[j-1][i],
                 jmp[j-1][i+(1<<(j-1))];
 }
 int index(int 1, int r) { // get index of min element
   int d = level(r-l+1);
   return comb(jmp[d][1], jmp[d][r-(1<<d)+1]);
 T query(int 1, int r) { return v[index(1,r)]; }
};
```

BIT.h

Description: N-D range sum query with point update

```
Time: \mathcal{O}\left((\log N)^D\right)
                                                        88cda4, 31 lines
template <class T, int ...Ns> struct BIT {
 T val = 0;
 void upd(T v) { val += v; }
 T query() { return val; }
template <class T, int N, int... Ns> struct BIT<T, N, Ns...> {
 BIT<T, Ns...> bit[N+1];
 template<typename... Args> void upd(int pos, Args... args) {
    for (; pos <= N; pos += (pos&-pos)) bit[pos].upd(args...);</pre>
 template<typename... Args> T sum(int r, Args... args) {
   T res = 0; for (; r; r -= (r\&-r)) res += bit[r].query(args
      \hookrightarrow . . . ) ;
    return res;
 template<typename... Args> T query(int 1, int r, Args... args
    →) {
    return sum(r,args...)-sum(l-1,args...);
}; // BIT<int,10,10> gives a 2D BIT
template < class T, int SZ> struct BITrange {
 BIT<T,SZ> bit[2]; // piecewise linear functions
 // let cum[x] = sum_{i=1}^{x}a[i]
 void upd(int hi, T val) { // add val to a[1..hi]
   bit[1].upd(1,val), bit[1].upd(hi+1,-val); // if x <= hi,
       \hookrightarrow cum[x] += val*x
   bit[0].upd(hi+1,hi*val); // if x > hi, cum[x] += val*hi
 void upd(int lo, int hi, T val) { upd(lo-1,-val), upd(hi,val)
     \hookrightarrow; }
```

SegTree.h

Description: 1D point update, range query **Time:** $\mathcal{O}(\log N)$

bf15d6, 21 lines

```
template<class T> struct Seg {
 const T ID = 0; // comb(ID,b) must equal b
 T comb(T a, T b) { return a+b; } // easily change this to min
    \hookrightarrow or max
  int n; vector<T> seq;
 void init(int _n) { n = _n; seg.rsz(2*n); }
  void pull(int p) { seg[p] = comb(seg[2*p], seg[2*p+1]); }
 void upd(int p, T value) { // set value at position p
   seq[p += n] = value;
   for (p /= 2; p; p /= 2) pull(p);
 T query(int 1, int r) { // sum on interval [1, r]
    T ra = ID, rb = ID; // make sure non-commutative operations
       \hookrightarrow work
    for (1 += n, r += n+1; 1 < r; 1 /= 2, r /= 2) {
     if (1&1) ra = comb(ra, seg[1++]);
      if (r\&1) rb = comb(seq[--r],rb);
    return comb(ra,rb);
};
```

PersSegTree.h

Description: persistent segtree with lazy updates, assumes that lazy[cur] is included in val[cur] before propagating cur ed6e9b, 60 lines

```
template < class T, int SZ> struct pseq {
 static const int LIMIT = 10000000; // adjust
 int l[LIMIT], r[LIMIT], nex = 0;
 T val[LIMIT], lazy[LIMIT];
 int copy(int cur) {
   int x = nex++;
   val[x] = val[cur], l[x] = l[cur], r[x] = r[cur], lazy[x] =
       ⇔lazv[cur];
    return x;
 T comb(T a, T b) { return min(a,b); }
 void pull(int x) { val[x] = comb(val[l[x]],val[r[x]]); }
 void push(int cur, int L, int R) {
   if (!lazy[cur]) return;
   if (L != R) {
     l[cur] = copy(l[cur]);
     val[l[cur]] += lazy[cur];
     lazy[l[cur]] += lazy[cur];
     r[cur] = copy(r[cur]);
     val[r[cur]] += lazy[cur];
     lazy[r[cur]] += lazy[cur];
   lazy[cur] = 0;
 T query(int cur, int lo, int hi, int L, int R) {
   if (lo <= L && R <= hi) return val[cur];</pre>
   if (R < lo || hi < L) return INF;
   int M = (L+R)/2;
```

```
return lazy[cur]+comb(query(1[cur],lo,hi,L,M), query(r[cur
     \hookrightarrow ], lo, hi, M+1, R));
int upd(int cur, int lo, int hi, T v, int L, int R) {
  if (R < lo || hi < L) return cur;
  int x = copy(cur);
  if (lo <= L && R <= hi) { val[x] += v, lazy[x] += v; return
     \hookrightarrow x; }
  push(x, L, R);
  int M = (L+R)/2;
  1[x] = upd(1[x], lo, hi, v, L, M), r[x] = upd(r[x], lo, hi, v, M+1, R)
     \hookrightarrow);
  pull(x); return x;
int build(vector<T>& arr, int L, int R) {
  int cur = nex++;
  if (L == R) {
   if (L < sz(arr)) val[cur] = arr[L];</pre>
   return cur;
  int M = (L+R)/2;
 l[cur] = build(arr, L, M), r[cur] = build(arr, M+1, R);
 pull(cur); return cur;
vi loc:
void upd(int lo, int hi, T v) { loc.pb(upd(loc.back(),lo,hi,v
   \hookrightarrow, 0, SZ-1)); }
T query(int ti, int lo, int hi) { return query(loc[ti],lo,hi
   \hookrightarrow , 0, SZ-1); }
void build(vector<T>& arr) { loc.pb(build(arr,0,SZ-1)); }
```

SegTreeBeats.h

Description: Interval min modifications

Time: ? 5688c0, 74 lines template<int SZ> struct SegTreeBeats { int N; 11 sum[2*SZ]: int mx[2][2*SZ], maxCnt[2*SZ]; void pull(int ind) { mx[0][ind] = max(mx[0][2*ind], mx[0][2*ind+1]);mx[1][ind] = max(mx[1][2*ind], mx[1][2*ind+1]);maxCnt[ind] = 0;FOR(i,2) { if $(mx[0][2*ind^i] == mx[0][ind])$ maxCnt[ind] += maxCnt \hookrightarrow [2*ind^i]; else $mx[1][ind] = max(mx[1][ind], mx[0][2*ind^i]);$ sum[ind] = sum[2*ind] + sum[2*ind+1];void build(vi& a, int ind = 1, int L = 0, int R = -1) { if (R == -1) R += N;if (L == R) { mx[0][ind] = sum[ind] = a[L];maxCnt[ind] = 1; mx[1][ind] = -1;return; int M = (L+R)/2; build(a, 2*ind, L, M); build(a, 2*ind+1, M+1, R); pull(ind);

```
void push (int ind, int L, int R)
   if (L == R) return;
    FOR(i,2)
      if (mx[0][2*ind^i] > mx[0][ind]) {
        sum[2*ind^i] -= (ll) maxCnt[2*ind^i] *
                 (mx[0][2*ind^i]-mx[0][ind]);
        mx[0][2*ind^i] = mx[0][ind];
 void upd(int x, int y, int t, int ind = 1, int L = 0, int R =
     \hookrightarrow -1) { // set a_i = min(a_i,t)
    if (R == -1) R += N;
    if (R < x || y < L || mx[0][ind] <= t) return;</pre>
    push (ind, L, R);
    if (x \le L \&\& R \le y \&\& mx[1][ind] \le t) {
      sum[ind] -= (11) maxCnt[ind] * (mx[0][ind]-t);
     mx[0][ind] = t;
     return:
    if (L == R) return;
    int M = (L+R)/2;
    upd(x,y,t,2*ind,L,M); upd(x,y,t,2*ind+1,M+1,R); pull(ind);
 11 qsum(int x, int y, int ind = 1, int L = 0, int R = -1) {
   if (R == -1) R += N;
    if (R < x \mid | y < L) return 0;
    push (ind, L, R);
    if (x <= L && R <= y) return sum[ind];
    int M = (L+R)/2;
    return qsum(x, y, 2*ind, L, M) + qsum(x, y, 2*ind+1, M+1, R);
 int qmax(int x, int y, int ind = 1, int L = 0, int R = -1) {
    if (R == -1) R += N;
    if (R < x \mid \mid y < L) return -1;
    push (ind, L, R);
    if (x <= L && R <= y) return mx[0][ind];
    int M = (L+R)/2;
    return max(qmax(x,y,2*ind,L,M), qmax(x,y,2*ind+1,M+1,R));
};
Treap.h
```

Description: Easiest BBST Time: $\mathcal{O}(\log N)$

typedef struct tnode* pt; struct tnode { int pri, val; pt c[2]; // essential int sz; 11 sum; // for range queries bool flip; // lazy update tnode (int _val) { pri = rand()+(rand()<<15); val = _val; c[0] = c[1] = NULL;</pre> sz = 1; sum = val;flip = 0;}; int getsz(pt x) { return x?x->sz:0; } 11 getsum(pt x) { return x?x->sum:0; }

```
pt prop(pt x) {
                if (!x || !x->flip) return x;
                swap (x->c[0], x->c[1]);
                x->flip = 0;
                FOR(i,2) if (x->c[i]) x->c[i]->flip ^= 1;
              pt calc(pt x) {
                assert(!x->flip);
                prop(x->c[0]), prop(x->c[1]);
                x->sz = 1+getsz(x->c[0])+getsz(x->c[1]);
                x\rightarrowsum = x\rightarrowval+getsum(x\rightarrowc[0])+getsum(x\rightarrowc[1]);
                return x;
              void tour(pt x, vi& v) {
                if (!x) return;
                prop(x):
                tour (x-c[0],v); v.pb (x-val); tour (x-c[1],v);
              pair<pt,pt> split(pt t, int v) { // >= v goes to the right
                if (!t) return {t,t};
                prop(t);
                if (t->val >= v) {
                  auto p = split(t->c[0], v); t->c[0] = p.s;
                  return {p.f, calc(t)};
                } else {
                  auto p = split(t->c[1], v); t->c[1] = p.f;
                  return {calc(t), p.s};
              pair<pt,pt> splitsz(pt t, int sz) { // leftmost sz nodes go to
                if (!t) return {t,t};
                prop(t);
                if (getsz(t->c[0]) >= sz) {
                  auto p = splitsz(t->c[0], sz); t->c[0] = p.s;
                  return {p.f, calc(t)};
                } else {
                  auto p = splitsz(t->c[1], sz-getsz(t->c[0])-1); t->c[1] = p
                  return {calc(t), p.s};
              pt merge(pt 1, pt r) {
                if (!1 || !r) return 1 ? 1 : r;
                prop(1), prop(r);
                if (1->pri > r->pri) 1->c[1] = merge(1->c[1],r), t = 1;
                else r - c[0] = merge(1, r - c[0]), t = r;
b45b6a, 77 lines
                return calc(t);
              pt ins(pt x, int v) { // insert v
                auto a = split(x,v), b = split(a.s,v+1);
                return merge(a.f, merge(new tnode(v),b.s));
              pt del(pt x, int v) { // delete v
                auto a = split(x,v), b = split(a.s,v+1);
                return merge(a.f,b.s);
```

Number Theory (3)

3.1 Modular Arithmetic

Modular.h

 $\textbf{Description:} \ \ \text{operations with modular arithmetic}$

20589d, 41 lines

```
template<class T> struct modular {
 T val:
  explicit operator T() const { return val; }
  modular() { val = 0; }
  modular(const 11& v) {
   val = (-MOD <= v && v <= MOD) ? v : v % MOD;
   if (val < 0) val += MOD;</pre>
  // friend ostream& operator<<(ostream& os, const modular& a)
    \hookrightarrow { return os << a.val; }
  friend void pr(const modular& a) { pr(a.val); }
  friend void re(modular& a) { ll x; re(x); a = modular(x); }
  friend bool operator == (const modular& a, const modular& b) {
    friend bool operator!=(const modular& a, const modular& b) {
    \hookrightarrowreturn ! (a == b); }
  friend bool operator<(const modular& a, const modular& b) {</pre>
    modular operator-() const { return modular(-val); }
  modular& operator+=(const modular& m) { if ((val += m.val) >=
    modular& operator-=(const modular& m) { if ((val -= m.val) <
    \hookrightarrow0) val += MOD; return *this; }
  modular& operator *= (const modular& m) { val = (11) val *m.val %
    →MOD; return *this; }
  friend modular pow(modular a, 11 p) {
   modular ans = 1; for (; p; p \neq 2, a \neq a) if (p\&1) ans \star=
    return ans;
  friend modular inv(const modular& a)
   assert (a != 0); return exp(a, MOD-2);
  modular& operator/=(const modular& m) { return (*this) *= inv
  friend modular operator+(modular a, const modular& b) {
    friend modular operator-(modular a, const modular& b) {
    friend modular operator* (modular a, const modular& b) {
    \hookrightarrowreturn a *= b; }
  friend modular operator/(modular a, const modular& b) {
    };
typedef modular<int> mi;
typedef pair<mi, mi> pmi;
typedef vector<mi> vmi;
typedef vector<pmi> vpmi;
```

ModInv.h

Description: pre-compute factorial mod inverses for MOD in linear time assume MOD is prime and SZ < MOD $_{\rm f88b07,\ 10\ lines}$

```
vl inv, fac, ifac;
void genInv(int SZ) {
```

```
inv.rsz(SZ), fac.rsz(SZ), ifac.rsz(SZ);
inv[1] = 1; FOR(i,2,SZ) inv[i] = MOD-MOD/i*inv[MOD%i]%MOD;
fac[0] = ifac[0] = 1;
FOR(i,1,SZ) {
   fac[i] = fac[i-1]*i%MOD;
   ifac[i] = ifac[i-1]*inv[i]%MOD;
}
```

ModMulLL.h

Description: multiply two 64-bit integers mod another if 128-bit is not available works for $0 \le a, b < mod < 2^{63}$

```
typedef unsigned long long ul;

// equivalent to (ul) (__int128(a) *b%mod)
ul modMul(ul a, ul b, const ul mod) {
    ll ret = a*b-mod*(ul) ((ld) a*b/mod);
    return ret+((ret<0)-(ret>=(ll) mod)) *mod;
}

ul modPow(ul a, ul b, const ul mod) {
    if (b == 0) return 1;
    ul res = modPow(a,b/2,mod);
    res = modMul(res,res,mod);
    if (b&1) return modMul(res,a,mod);
    return res;
}
```

ModSgrt.h

Description: find sqrt of integer mod a prime

a9a4c4, 26 lines

```
template<class T> T sqrt(modular<T> a) {
 auto p = pow(a, (MOD-1)/2); if (p != 1) return p == 0 ? 0:
    \hookrightarrow-1; // check if zero or does not have sgrt
 T s = MOD-1, e = 0; while (s % 2 == 0) s /= 2, e ++;
 modular < T > n = 1; while (pow(n, (MOD-1)/2) == 1) n = (T)(n)+1;
    \hookrightarrow // find non-square residue
 auto x = pow(a, (s+1)/2), b = pow(a, s), g = pow(n, s);
 int r = e;
 while (1) {
   auto B = b; int m = 0; while (B != 1) B \star= B, m ++;
   if (m == 0) return min((T)x, MOD-(T)x);
   FOR(i,r-m-1) g *= g;
   x *= g; g *= g; b *= g; r = m;
/* Explanation:
* Initially, x^2=ab, ord(b) = 2^m, ord(g) = 2^r where m < r
* g = g^{2^{r-m-1}} -> ord(g) = 2^{m+1}
* if x'=x*q, then b'=b*q^2
    (b')^{2^{m-1}} = (b*g^2)^{2^{m-1}}
             = b^{2^{m-1}} *q^{2^m}
             = -1 * -1
             = 7
 -> ord(b')|ord(b)/2
* m decreases by at least one each iteration
```

ModSum.h

 $\textbf{Description:} \ \ \text{Sums of mod'ed arithmetic progressions}$

typedef unsigned long long ul;
ul sumsq(ul to) { return (to-1)*to/2; } // sum of 0..to-1

3.2 Primality

```
PrimeSieve.h
```

 $\bf Description:$ tests primality up to n

Time: $\mathcal{O}(N \log \log N)$

5464fb, 13 lines

```
template<int SZ> struct Sieve {
  bitset<SZ> isprime;
  vi pr;

Sieve() {
   isprime.set(); isprime[0] = isprime[1] = 0;
   for (int i = 4; i < SZ; i += 2) isprime[i] = 0;
   for (int i = 3; i*i < SZ; i += 2) if (isprime[i])
     for (int j = i*i; j < SZ; j += i*2) isprime[j] = 0;
   vi pr;
  FOR(i,2,SZ) if (isprime[i]) pr.pb(i);
}
};</pre>
```

FactorFast.h

50ee96, 15 lines

Description: Factors integers up to 2⁶⁰

```
Time: ?
"PrimeSieve.h"
Sieve<1<<20> S = Sieve<1<<20>(); // should take care of all
   \hookrightarrow primes up to n^{(1/3)}
bool millerRabin(ll p) { // test primality
  if (p == 2) return true;
  if (p == 1 || p % 2 == 0) return false;
  11 s = p - 1; while (s % 2 == 0) s /= 2;
  FOR(i,30) { // strong liar with probability <= 1/4
    11 a = rand() % (p - 1) + 1, tmp = s;
    11 mod = mod_pow(a, tmp, p);
    while (tmp != p - 1 \&\& mod != 1 \&\& mod != p - 1) {
      mod = mod_mul(mod, mod, p);
      tmp *= 2;
    if (mod != p - 1 && tmp % 2 == 0) return false;
  return true;
ll f(ll a, ll n, ll &has) { return (mod_mul(a, a, n) + has) % n
vpl pollardsRho(ll d) {
  vpl res;
  auto& pr = S.pr;
  for (int i = 0; i < sz(pr) && pr[i]*pr[i] <= d; i++) if (d %
     \hookrightarrowpr[i] == 0) {
    int co = 0; while (d % pr[i] == 0) d /= pr[i], co ++;
    res.pb({pr[i],co});
```

Euclid CRT IntPerm MatroidIntersect PermGroup

3.3 Divisibility

Euclid.h

Description: Euclidean Algorithm

338527, 9 line

f295dd, 20 lines

CRT.h

Description: Chinese Remainder Theorem

Combinatorial (4)

IntPerm h

Time: $\mathcal{O}(n)$

Description: convert permutation of $\{0, 1, ..., n-1\}$ to integer in [0, n!) **Usage:** assert (encode (decode (5, 37)) == 37);

```
vi decode(int n, int a) {
  vi el(n), b; iota(all(el),0);
  FOR(i,n) {
    int z = a%sz(el);
    b.pb(el[z]); a /= sz(el);
    swap(el[z],el.back()); el.pop_back();
}
  return b;
}
int encode(vi b) {
```

```
int n = sz(b), a = 0, mul = 1;
  vi pos(n); iota(all(pos),0); vi el = pos;
    int z = pos[b[i]]; a += mul*z; mul *= sz(el);
    swap(pos[el[z]],pos[el.back()]);
    swap(el[z],el.back()); el.pop_back();
  return a:
MatroidIntersect.h
Description: max size of independent set in both graphic + colorful ma-
Time: ?
"DSU.h"
                                                    40170e, 108 lines
int R:
map<int, int> m;
struct Element {
 pi ed;
 int col;
  bool in independent set = 0;
  int independent_set_position;
  Element (int u, int v, int c) { ed = \{u,v\}; col = c; \}
vi independent set;
vector<Element> ground set;
bool col_used[300];
struct GBasis {
 DSU D:
  void reset() { D.init(sz(m)); }
  void add(pi v) { assert(D.unite(v.f,v.s)); }
 bool independent_with(pi v) { return !D.sameSet(v.f,v.s); }
};
GBasis basis, basis wo[300];
bool graph_oracle(int inserted) {
  return basis.independent_with(ground_set[inserted].ed);
bool graph_oracle(int inserted, int removed) {
  int wi = ground_set[removed].independent_set_position;
  return basis_wo[wi].independent_with(ground_set[inserted].ed)
void prepare_graph_oracle() {
  basis.reset();
  FOR(i,sz(independent_set)) basis_wo[i].reset();
  FOR(i,sz(independent_set)) {
    pi v = ground_set[independent_set[i]].ed; basis.add(v);
    FOR(j,sz(independent_set)) if (i != j) basis_wo[j].add(v);
bool colorful_oracle(int ins) {
  ins = ground set[ins].col;
  return !col_used[ins];
bool colorful_oracle(int ins, int rem) {
 ins = ground_set[ins].col;
  rem = ground set[rem].col;
  return !col_used[ins] || ins == rem;
void prepare_colorful_oracle() {
  FOR(i,R) col_used[i] = 0;
  trav(t,independent_set) col_used[ground_set[t].col] = 1;
```

```
bool augment() {
 prepare_graph_oracle();
 prepare_colorful_oracle();
 vi par(sz(ground_set),MOD);
 queue<int> q;
 FOR(i,sz(ground_set)) if (colorful_oracle(i)) {
   assert(!ground set[i].in independent set);
   par[i] = -1; q.push(i);
 int lst = -1;
 while (sz(q)) {
   int cur = q.front(); q.pop();
   if (ground_set[cur].in_independent_set) {
     FOR(to,sz(ground_set)) if (par[to] == MOD) {
       if (!colorful oracle(to,cur)) continue;
       par[to] = cur; q.push(to);
    } else {
      if (graph_oracle(cur)) { lst = cur; break; }
     trav(to,independent_set) if (par[to] == MOD) {
       if (!graph_oracle(cur,to)) continue;
       par[to] = cur; q.push(to);
 if (1st == -1) return 0;
   ground_set[lst].in_independent_set ^= 1;
   lst = par[lst];
 } while (lst !=-1);
 independent set.clear();
 FOR(i,sz(ground_set)) if (ground_set[i].in_independent_set) {
    ground_set[i].independent_set_position = sz(independent_set
    independent_set.pb(i);
 return 1;
void solve() {
 re(R); if (R == 0) exit(0);
 m.clear(); ground_set.clear(); independent_set.clear();
 FOR(i,R) {
   int a,b,c,d; re(a,b,c,d);
   ground_set.pb(Element(a,b,i));
   ground_set.pb(Element(c,d,i));
   m[a] = m[b] = m[c] = m[d] = 0;
 int co = 0;
 trav(t,m) t.s = co++;
 trav(t,ground_set) t.ed.f = m[t.ed.f], t.ed.s = m[t.ed.s];
 while (augment());
 ps(2*sz(independent_set));
```

PermGroup.h

Description: Schreier-Sims, count number of permutations in group and test whether permutation is a member of group

Time: ?

```
vi operator*(const vi& a, const vi& b) {
 vi c(sz(a)); FOR(i,sz(a)) c[i] = a[b[i]];
  return c;
struct Group {
 bool flag[N];
  vi sigma[N]; // sigma[t][k] = t, sigma[t][x] = x if x > k
  vector<vi> gen;
  void clear(int p) {
    memset (flag, 0, sizeof flag);
    flag[p] = 1; sigma[p] = id();
    gen.clear();
} q[N];
bool check(const vi& cur, int k) {
 if (!k) return 1;
  int t = cur[k];
  return q[k].flaq[t] ? check(inv(q[k].siqma[t])*cur,k-1) : 0;
void updateX(const vi& cur, int k);
void ins(const vi& cur, int k) {
 if (check(cur,k)) return;
  g[k].gen.pb(cur);
  FOR(i,n) if (g[k].flag[i]) updateX(cur*g[k].sigma[i],k);
void updateX(const vi& cur, int k) {
  int t = cur[k];
  if (q[k].flaq[t]) ins(inv(q[k].siqma[t])*cur,k-1); // fixes k
    \hookrightarrow -> k
  else {
   q[k].flaq[t] = 1, q[k].sigma[t] = cur;
    trav(x,g[k].gen) updateX(x*cur,k);
ll order (vector<vi> gen) {
  assert(sz(gen)); n = sz(gen[0]); FOR(i,n) g[i].clear(i);
  trav(a,gen) ins(a,n-1); // insert perms into group one by one
  11 \text{ tot} = 1;
  FOR(i,n) {
   int cnt = 0; FOR(j,i+1) cnt += g[i].flag[j];
   tot *= cnt;
  return tot;
```

Numerical (5)

5.1 Matrix

Matrix.h

Description: 2D matrix operations

c6abe5, 36 lines

```
FOR(i,r) FOR(j,c) d[i][j] += m.d[i][j];
    return *this;
 Mat& operator = (const Mat& m) {
   assert(r == m.r && c == m.c);
   FOR(i,r) FOR(j,c) d[i][j] -= m.d[i][j];
    return *this;
 Mat operator*(const Mat& m) {
    assert(c == m.r); Mat x(r,m.c);
    FOR(i,r) FOR(j,c) FOR(k,m.c) x.d[i][k] += d[i][j]*m.d[j][k]
       \hookrightarrow1:
    return x;
 Mat operator+(const Mat& m) { return Mat(*this)+=m; }
 Mat operator-(const Mat& m) { return Mat(*this)-=m; }
 Mat& operator*=(const Mat& m) { return *this = (*this)*m; }
 friend Mat pow(Mat m, 11 p) {
   assert (m.r == m.c);
   Mat r(m.r.m.c);
   FOR(i, m.r) r.d[i][i] = 1;
    for (; p; p /= 2, m \star= m) if (p&1) r \star= m;
    return r:
};
```

MatrixInv.h

Description: calculates determinant via gaussian elimination **Time:** $\mathcal{O}(N^3)$

```
00ad8c, 31 lines
template < class T > T gauss (Mat < T > & m) { // determinant of 1000
  \hookrightarrow x1000 Matrix in \sim1s
 int n = m.r;
 T prod = 1; int nex = 0;
 FOR(i,n) {
    int row = -1; // for 1d use EPS rather than 0
   FOR(j, nex, n) if (m.d[j][i] != 0) { row = j; break; }
    if (row == -1) { prod = 0; continue; }
    if (row != nex) prod *= -1, swap(m.d[row], m.d[nex]);
   prod *= m.d[nex][i];
    auto x = 1/m.d[nex][i]; FOR(k,i,m.c) m.d[nex][k] *= x;
   FOR(j,n) if (j != nex) {
     auto v = m.d[j][i];
     if (v != 0) FOR(k, i, m.c) m.d[j][k] -= v*m.d[nex][k];
   nex ++;
 return prod;
template < class T > Mat < T > inv (Mat < T > m) {
 int n = m.r;
 Mat < T > x(n, 2*n);
 FOR(i,n) {
   x.d[i][i+n] = 1;
   FOR(j,n) \times d[i][j] = m.d[i][j];
 if (gauss(x) == 0) return Mat < T > (0,0);
 Mat < T > r(n,n);
 FOR(i,n) FOR(j,n) r.d[i][j] = x.d[i][j+n];
 return r;
```

MatrixTree.h

Description: Kirchhoff's Matrix Tree Theorem: given adjacency matrix, calculates # of spanning trees $^{\text{cdb606, 13 lines}}$

```
mi numSpan(Mat<mi> m) {
   int n = m.r;
   Mat<mi> res(n-1,n-1);
   FOR(i,n) FOR(j,i+1,n) {
      mi ed = m.d[i][j];
      res.d[i][i] += ed;
      if (j != n-1) {
       res.d[j][j] += ed;
       res.d[i][j] -= ed, res.d[j][i] -= ed;
    }
   }
  return gauss(res);
}
```

5.2 Polynomials

Karatsuba.h

Description: multiply two polynomials **Time:** $\mathcal{O}\left(N^{\log_2 3}\right)$

```
21f372, 26 lines
int size(int s) { return s > 1 ? 32-_builtin_clz(s-1) : 0; }
void karatsuba(ll *a, ll *b, ll *c, ll *t, int n) {
 int ca = 0, cb = 0; FOR(i,n) ca += !!a[i], cb += !!b[i];
 if (min(ca, cb) <= 1500/n) { // few numbers to multiply
    if (ca > cb) swap(a, b);
    FOR(i,n) if (a[i]) FOR(j,n) c[i+j] += a[i]*b[j];
  } else {
    int h = n \gg 1;
    karatsuba(a, b, c, t, h); // a0*b0
    karatsuba(a+h, b+h, c+n, t, h); // a1*b1
    FOR(i,h) \ a[i] += a[i+h], \ b[i] += b[i+h];
    karatsuba(a, b, t, t+n, h); // (a0+a1) * (b0+b1)
    FOR(i,h) a[i] -= a[i+h], b[i] -= b[i+h];
    FOR(i,n) t[i] -= c[i] + c[i+n];
    FOR(i,n) c[i+h] += t[i], t[i] = 0;
vl conv(vl a, vl b) {
 int sa = sz(a), sb = sz(b); if (!sa | | !sb) return {};
  int n = 1 << size(max(sa,sb)); a.rsz(n), b.rsz(n);
  v1 c(2*n), t(2*n); FOR(i,2*n) t[i] = 0;
 karatsuba(&a[0], &b[0], &c[0], &t[0], n);
 c.rsz(sa+sb-1); return c;
```

FFT.h

Description: multiply two polynomials

Time: $O(N \log N)$

```
FOR(i,n) \ roots[i] = cd(cos(ang*i),sin(ang*i)); // is there a
     \hookrightarrow way to do this more quickly?
template<class T> void fft(vector<T>& a, const vector<T>& roots
  \hookrightarrow, bool inv = 0) {
  int n = sz(a);
  for (int i = 1, j = 0; i < n; i++) { // sort by reverse bit
     \hookrightarrowrepresentation
    int bit = n >> 1;
    for (; j&bit; bit >>= 1) j ^= bit;
    j ^= bit; if (i < j) swap(a[i], a[j]);</pre>
  for (int len = 2; len <= n; len <<= 1)
    for (int i = 0; i < n; i += len)
     FOR(i,len/2) {
        int ind = n/len*j; if (inv && ind) ind = n-ind;
        auto u = a[i+j], v = a[i+j+len/2] * roots[ind];
        a[i+j] = u+v, a[i+j+len/2] = u-v;
  if (inv) { T i = T(1)/T(n); trav(x,a) x *= i; }
template<class T> vector<T> mult(vector<T> a, vector<T> b) {
 int s = sz(a) + sz(b) - 1, n = 1 << size(s);
  vector<T> roots(n); genRoots(roots);
 a.rsz(n), fft(a,roots);
 b.rsz(n), fft(b,roots);
 FOR(i,n) a[i] \star = b[i];
  fft(a,roots,1); return a;
```

FFTmod.h

Description: multiply two polynomials with arbitrary MOD ensures precision by splitting in half

```
"FFT.h"
                                                           a8a6ed, 27 lines
vl multMod(const vl& a, const vl& b) {
  if (!min(sz(a),sz(b))) return {};
  int s = sz(a)+sz(b)-1, n = 1 << size(s), cut = sqrt(MOD);
  vcd roots(n); genRoots(roots);
  vcd ax(n), bx(n);
  FOR(i,sz(a)) ax[i] = cd((int)a[i]/cut, (int)a[i]%cut); // <math>ax(a)
     \hookrightarrow x) =a1 (x) +i *a0 (x)
  FOR(i,sz(b)) bx[i] = cd((int)b[i]/cut, (int)b[i]%cut); // bx(
     \hookrightarrow x) =b1 (x) +i *b0 (x)
  fft(ax,roots), fft(bx,roots);
  vcd v1(n), v0(n);
  FOR(i,n) {
    int j = (i ? (n-i) : i);
    v1[i] = (ax[i]+conj(ax[j]))*cd(0.5,0)*bx[i]; // v1 = a1*(b1)
       \hookrightarrow +b0 *cd(0,1));
    v0[i] = (ax[i]-conj(ax[j]))*cd(0,-0.5)*bx[i]; // v0 = a0*(
       \hookrightarrow b1+b0*cd(0,1));
  fft(v1, roots, 1), fft(v0, roots, 1);
  vl ret(n);
  F0R(i,n) {
    11 \ V2 = (11) \ round(v1[i].real()); // a1*b1
    11 V1 = (11) round(v1[i].imag())+(11) round(v0[i].real()); //
       \hookrightarrow a0*b1+a1*b0
    11 V0 = (11) round(v0[i].imag()); // a0*b0
    ret[i] = ((V2%MOD*cut+V1)%MOD*cut+V0)%MOD;
  ret.rsz(s); return ret;
} // \sim 0.8s when sz(a) = sz(b) = 1 << 19
```

```
PolyInv.h
Description: ?
Time: ?
                                                           a5fd4a, 11 lines
"PolyConv.h"
template<class T> vector<T> inv(vector<T> v, int p) { //
   \hookrightarrow compute inverse of v mod x^p, where v[0] = 1
  v.rsz(p); vector<T> a = {T(1)/v[0]};
  for (int i = 1; i < p; i *= 2) {
    if (2*i > p) v.rsz(2*i);
    auto l = vector < T > (begin(v), begin(v) + i), r = vector < T > (
        \hookrightarrow begin (v) +i, begin (v) +2*i);
    auto c = conv(a, 1); c = vector < T > (begin(c) + i, end(c));
    auto b = conv(a*T(-1), conv(a,r)+c); b.rsz(i);
    a.insert(end(a),all(b));
 a.rsz(p); return a;
PolvDiv.h
Description: divide two polynomials
Time: \mathcal{O}(N \log N)?
```

```
05bb2d, 7 lines
"PolvInv.h"
template<class T> pair<vector<T>, vector<T>> divi(const vector<T
   \hookrightarrow>& f, const vector<T>& g) { // f = q*q+r
  if (sz(f) < sz(g)) return {{},f};</pre>
 auto q = conv(inv(rev(q), sz(f) - sz(q) + 1), rev(f));
 q.rsz(sz(f)-sz(g)+1); q = rev(q);
 auto r = f-conv(q, g); r.rsz(sz(g)-1);
 return {q,r};
```

PolySart.h

Description: find sqrt of polynomial Time: $\mathcal{O}(N \log N)$?

```
"PolyInv.h"
                                                          784e58, 8 lines
template<class T> vector<T> sqrt(vector<T> v, int p) { // S*S =
  \hookrightarrow v mod x^p, p is power of 2
 assert (v[0] == 1); if (p == 1) return {1};
 v.rsz(p);
 auto S = sqrt(v, p/2);
 auto ans = S+conv(v,inv(S,p));
 ans.rsz(p); ans \star= T(1)/T(2);
 return ans:
```

5.3 Misc.

LinRec.h

Description: Berlekamp-Massey: computes linear recurrence of order n for sequence of 2n terms Time: ?

```
using namespace vecOp;
struct LinRec {
 vmi x; // original sequence
 vmi C, rC;
 void init(const vmi& _x) {
   x = _x; int n = sz(x), m = 0;
   vmi B; B = C = \{1\}; // B is fail vector
   mi b = 1; // B gives 0,0,0,...,b
   FOR(i,n) {
     m ++;
     mi d = x[i]; FOR(j,1,sz(C)) d += C[j]*x[i-j];
     if (d == 0) continue; // recurrence still works
```

```
auto _B = C; C.rsz(max(sz(C), m+sz(B)));
      mi coef = d/b; FOR(j, m, m+sz(B)) C[j] -= coef*B[j-m]; //
          \hookrightarrowrecurrence that gives 0,0,0,...,d
      if (sz(B) < m+sz(B)) \{ B = B; b = d; m = 0; \}
    rC = C; reverse(all(rC)); // polynomial for getPo
    C.erase(begin(C)); trav(t,C) t \star = -1; // x[i] = sum_{i=0}^{i=0} (sz)
       \hookrightarrow (C) -1}C[j] *x[i-j-1]
  vmi getPo(int n) {
    if (n == 0) return {1};
    vmi x = qetPo(n/2); x = rem(x*x,rC);
    if (n\&1) { vmi \ v = \{0,1\}; \ x = rem(x*v,rC); \}
    return x;
 mi eval(int n) {
    vmi t = getPo(n);
    mi ans = 0; FOR(i,sz(t)) ans += t[i]*x[i];
    return ans;
};
```

Integrate.h Description: ?

// db f(db x) { return x*x+3*x+1; }

```
db \quad quad(db \quad (*f) \quad (db), \quad db \quad a, \quad db \quad b) \quad \{
  const int n = 1000;
  db dif = (b-a)/2/n, tot = f(a)+f(b);
  FOR(i, 1, 2*n) tot += f(a+i*dif)*(i&1?4:2);
  return tot*dif/3;
```

IntegrateAdaptive.h

Description: ?

b48168, 19 lines

3ddcbc, 73 lines

693e87, 8 lines

```
// db f(db x) { return x*x+3*x+1; }
db simpson(db (*f)(db), db a, db b) {
 db c = (a+b) / 2;
 return (f(a) + 4*f(c) + f(b)) * (b-a) / 6;
db rec(db (*f)(db), db a, db b, db eps, db S) {
 db c = (a+b) / 2;
 db S1 = simpson(f, a, c);
  db S2 = simpson(f, c, b), T = S1 + S2;
 if (abs(T - S) \le 15 \times eps \mid | b-a < 1e-10)
   return T + (T - S) / 15;
  return rec(f, a, c, eps/2, S1) + rec(f, c, b, eps/2, S2);
db quad(db (*f)(db), db a, db b, db eps = 1e-8) {
 return rec(f, a, b, eps, simpson(f, a, b));
```

Simplex.h

49e624, 35 lines

Description: Simplex algorithm for linear programming, maximize $c^T x$ subject to Ax < b, x > 0

```
Time: ?
typedef double T;
```

```
typedef vector<T> vd;
typedef vector<vd> vvd;
const T eps = 1e-8, inf = 1/.0;
```

DSU ManhattanMST LCAjumps

```
#define ltj(X) if (s == -1 \mid | mp(X[j], N[j]) < mp(X[s], N[s])) s=
struct LPSolver {
 int m, n;
  vi N, B;
  vvd D;
  LPSolver(const vvd& A, const vd& b, const vd& c) :
   m(sz(b)), n(sz(c)), N(n+1), B(m), D(m+2), vd(n+2)) {
     FOR(i,m) FOR(j,n) D[i][j] = A[i][j];
     FOR(i,m) \{ B[i] = n+i; D[i][n] = -1; D[i][n+1] = b[i]; \}
         \hookrightarrow // B[i] -> basic variables, col n+1 is for constants
         \hookrightarrow, why D[i][n]=-1?
     FOR(j,n) \{ N[j] = j; D[m][j] = -c[j]; \} // N[j] -> non-
         ⇒basic variables, all zero
     N[n] = -1; D[m+1][n] = 1;
  void print() {
   ps("D");
   trav(t,D) ps(t);
   ps();
   ps("B",B);
   ps("N",N);
   ps();
  void pivot(int r, int s) { // row, column
   T \star a = D[r].data(), inv = 1/a[s]; // eliminate col s from
       \hookrightarrowconsideration
   FOR(i,m+2) if (i != r && abs(D[i][s]) > eps) {
     T *b = D[i].data(), inv2 = b[s]*inv;
     FOR(j,n+2) b[j] -= a[j]*inv2;
     b[s] = a[s] * inv2;
    FOR(j,n+2) if (j != s) D[r][j] *= inv;
   FOR(i, m+2) if (i != r) D[i][s] \star = -inv;
   D[r][s] = inv; swap(B[r], N[s]); // swap a basic and non-
      ⇒basic variable
  bool simplex(int phase) {
   int x = m+phase-1;
    for (;;) {
      int s = -1; FOR(j, n+1) if (N[j] != -phase) ltj(D[x]); //
         \hookrightarrow find most negative col
     if (D[x][s] >= -eps) return true; // have best solution
     int r = -1;
     FOR(i,m) {
       if (D[i][s] <= eps) continue;
       if (r == -1 \mid | mp(D[i][n+1] / D[i][s], B[i])
               < mp(D[r][n+1] / D[r][s], B[r])) r = i; // find
                  \hookrightarrowsmallest positive ratio
      if (r == -1) return false; // unbounded
     pivot(r, s);
  T solve(vd &x) {
   int r = 0; FOR(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i;
   if (D[r][n+1] < -eps) { // x=0 is not a solution}
     pivot(r, n); // -1 is artificial variable, initially set
         if (!simplex(2) || D[m+1][n+1] < -eps) return -inf; // no
         \hookrightarrow solution
```

Graphs (6)

6.1 Fundamentals

DSU.h

Description: ? Time: $\mathcal{O}(N\alpha(N))$

```
cbfb79, 22 lines
struct DSU {
 vi e;
 void init(int n) { e = vi(n, -1); }
 int qet(int x) { return e[x] < 0 ? x : e[x] = qet(e[x]); } //
    \hookrightarrow path compression
 bool sameSet(int a, int b) { return get(a) == get(b); }
 int size(int x) { return -e[get(x)]; }
 bool unite(int x, int y) { // union-by-rank
   x = get(x), y = get(y); if (x == y) return 0;
   if (e[x] > e[y]) swap(x,y);
   e[x] += e[y]; e[y] = x;
   return 1;
};
// computes the minimum spanning tree in O(ElogE) time
template<class T> T kruskal(int n, vector<pair<T,pi>> edge) {
 sort(all(edge));
 T ans = 0; DSU D; D.init(n);
 trav(a,edge) if (D.unite(a.s.f,a.s.s)) ans += a.f; // edge is
    \hookrightarrow in MST
 return ans;
```

ManhattanMST.h

Description: Compute MST of points where edges are manhattan distances **Time:** $\mathcal{O}\left(N\log N\right)$

```
"DSU.h" 6f801e, 62 lines
int N;
vector<array<int,3>> cur;
vector<pair<11,pi>> ed;
vi ind;

struct {
   map<int,pi> m;
   void upd(int a, pi b) {
      auto it = m.lb(a);
      if (it != m.end() && it->s <= b) return;
      m[a] = b; it = m.find(a);
   while (it != m.begin() && prev(it)->s >= b) m.erase(prev(it ->));
}
```

```
pi query(int y) { // for all a > y find min possible value of
    \hookrightarrow b
    auto it = m.ub(y);
    if (it == m.end()) return {2*MOD, 2*MOD};
    return it->s;
} S;
void solve() {
  sort(all(ind),[](int a, int b) { return cur[a][0] > cur[b
     \hookrightarrow1[0]; });
  S.m.clear();
  int nex = 0;
  trav(x, ind) { // cur[x][0] <= ?, cur[x][1] < ?}
    while (\text{nex} < N \&\& \text{cur[ind[nex]][0]} >= \text{cur[x][0]}) 
      int b = ind[nex++];
      S.upd(cur[b][1], {cur[b][2],b});
    pi t = S.query(cur[x][1]);
    if (t.s != 2*MOD) ed.pb({(11)t.f-cur[x][2],{x,t.s}});
ll mst(vpi v) {
  N = sz(v); cur.resz(N); ed.clear();
  ind.clear(); FOR(i,N) ind.pb(i);
  sort(all(ind),[&v](int a, int b) { return v[a] < v[b]; });</pre>
  FOR(i, N-1) if (v[ind[i]] == v[ind[i+1]]) ed.pb(\{0, \{ind[i], ind\}\})
     \hookrightarrow[i+1]}});
  FOR(i,2) { // it's probably ok to consider just two guadrants
     \hookrightarrow ?
    FOR(i, N) {
      auto a = v[i];
      cur[i][2] = a.f+a.s;
    FOR(i,N) { // first octant
      auto a = v[i];
      cur[i][0] = a.f-a.s;
      cur[i][1] = a.s;
    solve();
    FOR(i,N) { // second octant
      auto a = v[i];
      cur[i][0] = a.f;
      cur[i][1] = a.s-a.f;
    solve():
    trav(a,v) a = {a.s,-a.f}; // rotate 90 degrees, repeat
  return kruskal (ed);
```

6.2 Trees

LCAjumps.h

Description: calculates least common ancestor in tree with binary jumping **Time:** $\mathcal{O}\left(N\log N\right)$

```
template<int SZ> struct LCA {
   static const int BITS = 32__builtin_clz(SZ);
   int N, R = 1; // vertices from 1 to N, R = root
   vi adj[SZ];
   int par[BITS][SZ], depth[SZ];

// INITIALIZE
   void addEdge(int u, int v) { adj[u].pb(v), adj[v].pb(u); }
```

```
void dfs(int u, int prev){
   par[0][u] = prev;
    depth[u] = depth[prev]+1;
   trav(v,adj[u]) if (v != prev) dfs(v, u);
  void init(int _N) {
   N = N; dfs(R, 0);
   FOR(k, 1, BITS) FOR(i, 1, N+1) par[k][i] = par[k-1][par[k-1][i]
  // OUERY
  int getPar(int a, int b) {
   ROF(k,BITS) if (b&(1<< k)) a = par[k][a];
    return a;
  int lca(int u, int v){
   if (depth[u] < depth[v]) swap(u,v);</pre>
   u = getPar(u,depth[u]-depth[v]);
   ROF(k,BITS) if (par[k][u] != par[k][v]) u = par[k][u], v =
      \rightarrowpar[k][v];
    return u == v ? u : par[0][u];
  int dist(int u, int v) {
    return depth[u]+depth[v]-2*depth[lca(u,v)];
};
```

CentroidDecomp.h

Description: can support tree path queries and updates

```
Time: \mathcal{O}(N \log N)
                                                      81e9e4, 45 lines
template<int SZ> struct CD {
 vi adj[SZ];
  bool done[SZ]:
  int sub[SZ], par[SZ];
 vl dist[SZ];
  pi cen[SZ];
  void addEdge(int a, int b) { adj[a].pb(b), adj[b].pb(a); }
  void dfs (int x) {
    sub[x] = 1;
   trav(y,adj[x]) if (!done[y] && y != par[x]) {
     par[y] = x; dfs(y);
     sub[x] += sub[y];
  int centroid(int x) {
   par[x] = -1; dfs(x);
   for (int sz = sub[x];;) {
     pi mx = \{0, 0\};
     trav(y,adj[x]) if (!done[y] && y != par[x])
       ckmax(mx, {sub[y],y});
     if (mx.f*2 <= sz) return x;
     x = mx.s;
  void genDist(int x, int p) {
   dist[x].pb(dist[p].back()+1);
   trav(y,adj[x]) if (!done[y] \&\& y != p) {
     cen[y] = cen[x];
     genDist(y,x);
  void gen(int x, bool fst = 0) {
   done[x = centroid(x)] = 1; dist[x].pb(0);
   if (fst) cen[x].f = -1;
   int co = 0;
```

```
trav(y,adj[x]) if (!done[y]) {
    cen[y] = {x, co++};
    genDist(y,x);
  trav(y,adj[x]) if (!done[y]) gen(y);
void init() { gen(1,1); }
```

```
HLD.h
Description: Heavy Light Decomposition
Time: \mathcal{O}(\log^2 N) per path operations
                                                      69f40a, 50 lines
template<int SZ, bool VALUES_IN_EDGES> struct HLD {
 int N; vi adj[SZ];
 int par[SZ], sz[SZ], depth[SZ];
  int root[SZ], pos[SZ];
  LazySegTree<11,SZ> tree;
  void addEdge(int a, int b) { adj[a].pb(b), adj[b].pb(a); }
  void dfs sz(int v = 1) {
    if (par[v]) adj[v].erase(find(all(adj[v]),par[v]));
    sz[v] = 1;
    trav(u,adj[v]) {
      par[u] = v; depth[u] = depth[v]+1;
      dfs_sz(u); sz[v] += sz[u];
      if (sz[u] > sz[adj[v][0]]) swap(u, adj[v][0]);
  void dfs hld(int v = 1) {
    static int t = 0;
    pos[v] = t++;
    trav(u,adi[v]) {
     root[u] = (u == adj[v][0] ? root[v] : u);
      dfs hld(u);
  void init(int N) {
    N = N; par[1] = depth[1] = 0; root[1] = 1;
    dfs sz(); dfs hld();
  template <class BinaryOperation>
  void processPath(int u, int v, BinaryOperation op) {
    for (; root[u] != root[v]; v = par[root[v]]) {
      if (depth[root[u]] > depth[root[v]]) swap(u, v);
      op(pos[root[v]], pos[v]);
    if (depth[u] > depth[v]) swap(u, v);
    op(pos[u]+VALUES_IN_EDGES, pos[v]);
  void modifyPath(int u, int v, int val) { // add val to

→vertices/edges along path

    processPath(u, v, [this, &val](int 1, int r) { tree.upd(1,
       \hookrightarrowr, val); });
  void modifySubtree(int v, int val) { // add val to vertices/
     ⇒edges in subtree
    tree.upd(pos[v]+VALUES_IN_EDGES,pos[v]+sz[v]-1,val);
 11 queryPath(int u, int v) { // query sum of path
   11 res = 0; processPath(u, v, [this, &res](int 1, int r) {
       \hookrightarrow res += tree.qsum(1, r); });
    return res;
};
```

6.3 DFS Algorithms

SCC.h

Description: Kosaraju's Algorithm does DFS two times to generate SCC in topological order Time: $\mathcal{O}(N+M)$

9

f53f41, 24 lines

```
template<int SZ> struct SCC {
 int N, comp[SZ];
 vi adj[SZ], radj[SZ], todo, allComp;
 bitset<SZ> visit;
 void addEdge(int a, int b) { adj[a].pb(b), radj[b].pb(a); }
 void dfs(int v) {
   visit[v] = 1;
   trav(w,adj[v]) if (!visit[w]) dfs(w);
   todo.pb(v);
 void dfs2(int v, int val) {
   comp[v] = val;
   trav(w, radj[v]) if (comp[w] == -1) dfs2(w, val);
 void init(int _N) { // fills allComp
   N = N;
   FOR(i,N) comp[i] = -1, visit[i] = 0;
   FOR(i,N) if (!visit[i]) dfs(i);
   reverse(all(todo)); // now todo stores vertices in order of

→ topological sort

   trav(i,todo) if (comp[i] == -1) dfs2(i,i), allComp.pb(i);
```

2SAT.h

};

Description: ?

```
"SCC.h"
                                                     6c209d, 38 lines
template<int SZ> struct TwoSat {
 SCC<2*SZ> S;
 bitset<SZ> ans;
 int N = 0:
 int addVar() { return N++; }
 void either(int x, int y) {
   x = max(2*x, -1-2*x), y = max(2*y, -1-2*y);
   S.addEdge(x^1, y); S.addEdge(y^1, x);
 void implies(int x, int y) { either(\sim x, y); }
 void setVal(int x) { either(x,x); }
 void atMostOne(const vi& li) {
   if (sz(li) <= 1) return;
   int cur = \simli[0];
    FOR(i,2,sz(li)) {
     int next = addVar();
     either(cur,~li[i]);
     either(cur,next);
     either(~li[i],next);
     cur = ~next;
    either(cur,~li[1]);
 bool solve(int _N) {
   if (_N != -1) N = _N;
   S.init(2*N);
   for (int i = 0; i < 2*N; i += 2)
     if (S.comp[i] == S.comp[i^1]) return 0;
    reverse(all(S.allComp));
    vi tmp(2*N);
```

```
trav(i, S.allComp) if (tmp[i] == 0)
     tmp[i] = 1, tmp[S.comp[i^1]] = -1;
   FOR(i,N) if (tmp[S.comp[2*i]] == 1) ans[i] = 1;
   return 1;
};
```

EulerPath.h

Description: Eulerian Path for both directed and undirected graphs Time: $\mathcal{O}(N+M)$

fd7ad7, 30 lines template<int SZ, bool directed> struct Euler { int N. M = 0: vpi adj[SZ]; vpi::iterator its[SZ]; vector<bool> used; void addEdge(int a, int b) { if (directed) adj[a].pb({b,M}); else adj[a].pb({b,M}), adj[b].pb({a,M}); used.pb(0); M ++;vpi solve(int N, int src = 1) { N = N;FOR(i,1,N+1) its[i] = begin(adi[i]); vector<pair<pi,int>> ret, $s = \{\{\{src, -1\}, -1\}\};$ while (sz(s)) { int x = s.back().f.f;auto& it = its[x], end = adj[x].end(); while (it != end && used[it->s]) it ++; if (it == end) { if (sz(ret) && ret.back().f.s != s.back().f.f) return \hookrightarrow {}: // path isn't valid ret.pb(s.back()), s.pop back(); } else { s.pb($\{\{it->f,x\},it->s\}$); used[it->s] = 1; } if (sz(ret) != M+1) return {}; vpi ans; trav(t,ret) ans.pb({t.f.f,t.s}); reverse(all(ans)); return ans;

BCC.h

};

Description: computes biconnected components

Time: $\mathcal{O}(N+M)$

```
393aff, 37 lines
template<int SZ> struct BCC {
 int N:
  vpi adj[SZ], ed;
  void addEdge(int u, int v) {
   adj[u].pb({v,sz(ed)}), adj[v].pb({u,sz(ed)});
   ed.pb({u,v});
  int disc[SZ];
 vi st; vector<vi> fin;
  int bcc(int u, int p = -1) { // return lowest disc
   static int ti = 0;
   disc[u] = ++ti; int low = disc[u];
   int child = 0;
   trav(i,adj[u]) if (i.s != p)
     if (!disc[i.f]) {
       child ++; st.pb(i.s);
       int LOW = bcc(i.f,i.s); ckmin(low,LOW);
        // disc[u] < LOW -> bridge
       if (disc[u] <= LOW) {
         // if (p != -1 || child > 1) -> u is articulation
             \hookrightarrowpoint.
```

```
vi tmp; while (st.back() != i.s) tmp.pb(st.back()),
             →st.pop_back();
          tmp.pb(st.back()), st.pop_back();
          fin.pb(tmp);
      } else if (disc[i.f] < disc[u]) {</pre>
        ckmin(low,disc[i.f]);
        st.pb(i.s);
    return low;
  void init(int N) {
   N = N; FOR(i, N) disc[i] = 0;
    FOR(i,N) if (!disc[i]) bcc(i); // st should be empty after
       \hookrightarroweach iteration
};
```

6.4 Flows

```
Dinic.h
```

Description: faster flow

Time: $\mathcal{O}(N^2M)$ flow, $\mathcal{O}(M\sqrt{N})$ bipartite matching

f1366f, 47 lines

```
template<int SZ> struct Dinic {
 typedef 11 F; // flow type
 struct Edge { int to, rev; F f, c; };
 int N,s,t;
 vector<Edge> adj[SZ];
 typename vector<Edge>::iterator cur[SZ];
 void addEdge(int u, int v, F cap) {
   assert(cap >= 0); // don't try smth dumb
   Edge a\{v, sz(adj[v]), 0, cap\}, b\{u, sz(adj[u]), 0, 0\};
   adi[u].pb(a), adi[v].pb(b);
 int level[SZl:
 bool bfs() { // level = shortest distance from source
   // after computing flow, edges {u,v} such that level[u] \
      \hookrightarrowneg -1, level[v] = -1 are part of min cut
   FOR(i, N) level[i] = -1, cur[i] = begin(adj[i]);
   queue < int > q({s}); level[s] = 0;
   while (sz(q)) {
     int u = q.front(); q.pop();
     trav(e,adj[u]) if (level[e.to] < 0 && e.f < e.c) {
       level[e.to] = level[u]+1; q.push(e.to);
   return level[t] >= 0;
 F sendFlow(int v, F flow) {
   if (v == t) return flow;
   for (; cur[v] != end(adj[v]); cur[v]++) {
     Edge& e = *cur[v];
     if (level[e.to] != level[v]+1 || e.f == e.c) continue;
     auto df = sendFlow(e.to,min(flow,e.c-e.f));
     if (df) { // saturated at least one edge
       e.f += df; adj[e.to][e.rev].f -= df;
       return df:
   return 0;
 F maxFlow(int _N, int _s, int _t) {
   N = N, s = s, t = t; if (s == t) return -1;
```

```
F tot = 0;
    while (bfs()) while (auto flow = sendFlow(s, numeric limits<
       \hookrightarrowF>::max())) tot += flow;
    return tot;
};
```

```
MCMF.h
Description: Min-Cost Max Flow, no negative cycles allowed
Time: ?
                                                     f67674, 56 lines
template<class T> using pqg = priority_queue<T,vector<T>,
  \hookrightarrowgreater<T>>;
template<class T> T poll(pqq<T>& x) {
 T y = x.top(); x.pop();
 return y;
template<int SZ> struct mcmf {
 struct Edge { int to, rev; ll f, c, cost; };
 vector<Edge> adi[SZ];
 void addEdge(int u, int v, ll cap, ll cost) {
   assert(cap >= 0);
    Edge a\{v, sz(adj[v]), 0, cap, cost\}, b\{u, sz(adj[u]), 0, 0,
      \hookrightarrow -cost};
    adj[u].pb(a), adj[v].pb(b);
 int N, s, t;
 pi pre[SZ]: // previous vertex, edge label on path
 pl cost[SZ]; // tot cost of path, amount of flow
  11 totFlow, totCost, curCost;
  void reweight() { // ensures all non-negative edge weights
    FOR(i,N) trav(p,adj[i]) p.cost += cost[i].f-cost[p.to].f;
 bool spfa() { // reweighting will ensure that there will be
    --- negative weights only during the first time you run this
    FOR(i,N) cost[i] = {INF,0};
    cost[s] = \{0, INF\};
    pqg<pair<11, int>> todo({{0,s}});
    while (sz(todo)) {
      auto x = poll(todo); if (x.f > cost[x.s].f) continue;
      trav(a,adi[x,s]) if (x,f+a,cost < cost[a,to],f && a,f < a
        →.c) {
        pre[a.to] = {x.s,a.rev};
        cost[a.to] = {x.f+a.cost, min(a.c-a.f,cost[x.s].s)};
        todo.push({cost[a.to].f,a.to});
    curCost += cost[t].f; return cost[t].s;
 void backtrack() {
    auto f = cost[t].s; totFlow += f, totCost += curCost*f;
    for (int x = t; x != s; x = pre[x].f) {
     adj[x][pre[x].s].f -= f;
      adj[pre[x].f][adj[x][pre[x].s].rev].f += f;
 pl calc(int _N, int _s, int _t) {
    N = N; s = s, t = t; totFlow = totCost = curCost = 0;
    spfa();
    while (1)
     reweight();
     if (!spfa()) return {totFlow, totCost};
      backtrack();
```

```
};
```

};

GomoryHu.h

Description: Compute max flow between every pair of vertices of undirected

fe44db, 56 lines

```
template<int SZ> struct GomoryHu {
 int N;
  vector<pair<pi,int>> ed;
  void addEdge(int a, int b, int c) { ed.pb({{a,b},c}); }
  vector<vi> cor = {{}}; // groups of vertices
  map<int,int> adj[2*SZ]; // current edges of tree
  int side[SZ];
  int gen(vector<vi> cc) {
   Dinic<SZ> D = Dinic<SZ>();
   vi comp(N+1); FOR(i,sz(cc)) trav(t,cc[i]) comp[t] = i;
   trav(t,ed) if (comp[t.f.f] != comp[t.f.s]) {
     D.addEdge(comp[t.f.f],comp[t.f.s],t.s);
     D.addEdge(comp[t.f.s],comp[t.f.f],t.s);
    int f = D.maxFlow(0,1);
   FOR(i, sz(cc)) trav(j, cc[i]) side[j] = D.level[i] >= 0; //
      ⇒min cut
    return f:
  void fill(vi& v, int a, int b) {
   trav(t,cor[a]) v.pb(t);
   trav(t,adj[a]) if (t.f != b) fill (v,t.f,a);
  void addTree(int a, int b, int c) { adj[a][b] = c, adj[b][a]
  void delTree(int a, int b) { adj[a].erase(b), adj[b].erase(a)
  vector<pair<pi,int>> init(int _N) { // returns edges of
    \hookrightarrow Gomory-Hu Tree
   N = N;
   FOR(i,1,N+1) cor[0].pb(i);
    queue<int> todo; todo.push(0);
    while (sz(todo)) {
     int x = todo.front(); todo.pop();
     vector<vi> cc; trav(t,cor[x]) cc.pb({t});
     trav(t,adj[x]) {
       cc.pb({});
        fill(cc.back(),t.f,x);
      int f = gen(cc); // run max flow
     cor.pb({}), cor.pb({});
     trav(t, cor[x]) cor[sz(cor)-2+side[t]].pb(t);
     FOR(i,2) if (sz(cor[sz(cor)-2+i]) > 1) todo.push(sz(cor)
     FOR(i,sz(cor)-2) if (i != x \&\& adj[i].count(x)) {
       addTree(i,sz(cor)-2+side[cor[i][0]],adj[i][x]);
       delTree(i,x);
     } // modify tree edges
     addTree(sz(cor)-2,sz(cor)-1,f);
    vector<pair<pi,int>> ans;
   FOR(i,sz(cor)) trav(j,adj[i]) if (i < j.f)
     ans.pb({{cor[i][0],cor[j.f][0]},j.s});
    return ans;
```

6.5 Matching

DFSmatch.h

Description: naive bipartite matching Time: $\mathcal{O}(NM)$

37ad8b, 26 lines template<int SZ> struct MaxMatch { int N, flow = 0, match[SZ], rmatch[SZ]; bitset<SZ> vis; vi adj[SZ]; MaxMatch() { memset(match, 0, sizeof match); memset(rmatch, 0, sizeof rmatch); void connect(int a, int b, bool c = 1) { if (c) match[a] = b, rmatch[b] = a; else match[a] = rmatch[b] = 0; bool dfs(int x) { if (!x) return 1; if (vis[x]) return 0; vis[x] = 1;trav(t,adj[x]) if (t != match[x] && dfs(rmatch[t])) return connect(x,t),1; return O: void tri(int x) { vis.reset(); flow += dfs(x); } void init(int _N) { N = N; FOR(i, 1, N+1) if (!match[i]) tri(i);

Hungarian.h

};

Description: finds min cost to complete n jobs w/ m workers each worker is assigned to at most one job $(n \le m)$

```
Time: ?
int HungarianMatch (const vector<vi>& a) { // cost array,
  \hookrightarrownegative values are ok
 int n = sz(a)-1, m = sz(a[0])-1; // jobs 1..., workers 1...m
 vi u(n+1), v(m+1), p(m+1); //p[j] \rightarrow job\ picked\ by\ worker\ j
 FOR(i,1,n+1) { // find alternating path with job i
   p[0] = i; int j0 = 0;
    vi dist(m+1, MOD), pre(m+1,-1); // dist, previous vertex on

→ shortest path

    vector<bool> done(m+1, false);
      done[j0] = true;
      int i0 = p[j0], j1; int delta = MOD;
      FOR(j,1,m+1) if (!done[j]) {
       auto cur = a[i0][j]-u[i0]-v[j];
        if (cur < dist[j]) dist[j] = cur, pre[j] = j0;</pre>
        if (dist[j] < delta) delta = dist[j], j1 = j;</pre>
      FOR(j,m+1) // just dijkstra with potentials
        if (done[j]) u[p[j]] += delta, v[j] -= delta;
        else dist[j] -= delta;
      j0 = j1;
    } while (p[j0]);
    do { // update values on alternating path
     int j1 = pre[j0];
     p[j0] = p[j1];
      j0 = j1;
   } while (j0);
 return -v[0]; // min cost
```

UnweightedMatch.h

return false;

Description: general unweighted matching

Time: ? c24787, 79 lines template<int SZ> struct UnweightedMatch { int vis[SZ], par[SZ], orig[SZ], match[SZ], aux[SZ], t, N; // \hookrightarrow 1-based index vi adj[SZ]; queue<int> 0; void addEdge(int u, int v) { adj[u].pb(v); adj[v].pb(u); void init(int n) { N = n; t = 0;FOR(i,N+1) { adj[i].clear(); match[i] = aux[i] = par[i] = 0;void augment(int u, int v) { int pv = v, nv; pv = par[v]; nv = match[pv]; match[v] = pv; match[pv] = v; v = nv;} while(u != pv); int lca(int v, int w) { while (1) { if (v) { if (aux[v] == t) return v; aux[v] = t; v = orig[par[match[v]]]; swap(v, w); void blossom(int v, int w, int a) { while (orig[v] != a) { par[v] = w; w = match[v];if (vis[w] == 1) O.push(w), vis[w] = 0;orig[v] = orig[w] = a;v = par[w];bool bfs(int u) { fill(vis+1, vis+1+N, -1); iota(orig + 1, orig + N + 1, 1);Q = queue < int > (); Q.push(u); vis[u] = 0;while (sz(Q)) { int v = Q.front(); Q.pop(); trav(x,adi[v]) { if (vis[x] == -1) { par[x] = v; vis[x] = 1;if (!match[x]) return augment(u, x), true; Q.push(match[x]); vis[match[x]] = 0; } else if (vis[x] == 0 && orig[v] != orig[x]) { int a = lca(orig[v], orig[x]); blossom(x, v, a); blossom(v, x, a);

MaximalCliques LCT DirectedMST

```
int match() {
   int ans = 0;
    // find random matching (not necessary, constant
      →improvement)
   vi V(N-1); iota(all(V), 1);
    shuffle(all(V), mt19937(0x94949));
   trav(x, V) if(!match[x])
     trav(y,adj[x]) if (!match[y]) {
       match[x] = y, match[y] = x;
        ++ans; break;
   FOR(i,1,N+1) if (!match[i] && bfs(i)) ++ans;
};
```

\mathbf{Misc}

MaximalCliques.h

Description: Finds all maximal cliques

```
Time: \mathcal{O}\left(3^{n/3}\right)
```

```
f70515, 19 lines
typedef bitset<128> B;
int N:
B adj[128];
void cliques (B P = \simB(), B X={}, B R={}) { // possibly in
   ⇒clique, not in clique, in clique
  if (!P.any()) {
    if (!X.any()) {
     // do smth with maximal clique
   return:
  auto q = (P|X)._Find_first();
  auto cands = P&~eds[q]; // clique must contain q or non-
    ⇔neighbor of g
  FOR(i,N) if (cands[i]) {
   R[i] = 1;
    cliques(eds, f, P & eds[i], X & eds[i], R);
   R[i] = P[i] = 0; X[i] = 1;
```

LCT.h

Description: Link-Cut Tree, use vir for subtree size queries Time: $\mathcal{O}(\log N)$

06a240, 96 lines

```
typedef struct snode* sn;
struct snode {
  sn p, c[2]; // parent, children
  int val; // value in node
  int sum, mn, mx; // sum of values in subtree, min and max
    \hookrightarrowprefix sum
  bool flip = 0;
  // int vir = 0; stores sum of virtual children
  snode(int v) {
   p = c[0] = c[1] = NULL;
   val = v; calc();
  friend int getSum(sn x) { return x?x->sum:0; }
  friend int getMn(sn x) { return x?x->mn:0; }
  friend int getMx(sn x) { return x?x->mx:0; }
```

```
void prop() {
  if (!flip) return;
  swap(c[0],c[1]); tie(mn,mx) = mp(sum-mx,sum-mn);
  FOR(i,2) if (c[i]) c[i]->flip ^= 1;
  flip = 0:
void calc() {
  FOR(i,2) if (c[i]) c[i]->prop();
  int s0 = getSum(c[0]), s1 = getSum(c[1]); sum = s0+val+s1;
  mn = min(getMn(c[0]), s0+val+getMn(c[1]));
  mx = max(getMx(c[0]), s0+val+getMx(c[1]));
int dir() {
  if (!p) return -2;
  FOR(i,2) if (p\rightarrow c[i] == this) return i;
  return -1; // p is path-parent pointer, not in current
     ⇒splav tree
bool isRoot() { return dir() < 0; }</pre>
friend void setLink(sn x, sn y, int d) {
  if (y) y->p = x;
  if (d >= 0) x -> c[d] = v;
void rot() { // assume p and p->p propagated
  assert(!isRoot()); int x = dir(); sn pa = p;
  setLink(pa->p, this, pa->dir());
  setLink(pa, c[x^1], x);
  setLink(this, pa, x^1);
  pa->calc(); calc();
void splay() {
  while (!isRoot() && !p->isRoot()) {
    p->p->prop(), p->prop(), prop();
    dir() == p->dir() ? p->rot() : rot();
    rot();
  if (!isRoot()) p->prop(), prop(), rot();
  prop();
void access() { // bring this to top of tree
  for (sn v = this, pre = NULL; v; v = v -> p) {
    v->splay();
    // if (pre) v->vir -= pre->sz;
    // if (v->c[1]) v->vir += v->c[1]->sz;
    v->c[1] = pre; v->calc();
    pre = v;
    // v->sz should remain the same if using vir
  splay(); assert(!c[1]); // left subtree of this is now path

→ to root, right subtree is empty
void makeRoot() { access(); flip ^= 1; }
void set(int v) { splay(); val = v; calc(); } // change value
   \hookrightarrow in node, splay suffices instead of access because it
   ⇒doesn't affect values in nodes above it
friend sn lca(sn x, sn y) {
  if (x == y) return x;
  x->access(), y->access(); if (!x->p) return NULL; // access
     \hookrightarrow at y did not affect x, so they must not be connected
  x\rightarrow splay(); return x\rightarrow p ? x\rightarrow p : x;
friend bool connected(sn x, sn y) { return lca(x,y); }
friend int balanced(sn x, sn y) {
```

```
x->makeRoot(); y->access();
    return y->sum-2*y->mn;
  friend bool link(sn x, sn y) { // make x parent of y
    if (connected(x,y)) return 0; // don't induce cycle
    y->makeRoot(); y->p = x;
    // x->access(); x->sz += y->sz; x->vir += y->sz;
    return 1; // success!
  friend bool cut(sn x, sn y) { // x is originally parent of y
    x->makeRoot(); y->access();
    if (y->c[0] != x || x->c[0] || x->c[1]) return 0; // splay
       \hookrightarrowtree with y should not contain anything else besides x
    x\rightarrow p = y\rightarrow c[0] = NULL; y\rightarrow calc(); return 1; // calc is

→ redundant as it will be called elsewhere anyways?
};
```

12

DirectedMST.h

Description: computes the minimum directed spanning tree

Time: $\mathcal{O}(M \log M)$ 8fe6d9, 47 lines

```
struct Edge { int a, b; ll w; };
struct Node {
 Edge kev;
 Node *1, *r;
 ll delta;
 void prop()
   kev.w += delta;
   if (1) 1->delta += delta;
   if (r) r->delta += delta;
    delta = 0;
 Edge top() { prop(); return key; }
Node *merge(Node *a, Node *b) {
 if (!a || !b) return a ?: b;
 a->prop(), b->prop();
 if (a->key.w > b->key.w) swap(a, b);
 swap(a->1, (a->r = merge(b, a->r)));
 return a;
void pop(Node*\& a) { a->prop(); a = merge(a->1, a->r); }
11 dmst(int n, int r, vector<Edge>& g) {
 DSU dsu: dsu.init(n):
 vector<Node*> heap(n);
 trav(e, g) heap[e.b] = merge(heap[e.b], new Node{e});
 11 \text{ res} = 0;
 vi seen(n, -1), path(n); seen[r] = r;
 FOR(s,n) {
    int u = s, qi = 0, w;
    while (seen[u] < 0) {</pre>
      path[qi++] = u, seen[u] = s;
      if (!heap[u]) return -1;
      Edge e = heap[u]->top();
      heap[u]->delta -= e.w, pop(heap[u]);
      res += e.w, u = dsu.get(e.a);
      if (seen[u] == s) {
        Node * cyc = 0;
        do cyc = merge(cyc, heap[w = path[--qi]]);
        while (dsu.unite(u, w));
        u = dsu.get(u);
        heap[u] = cyc, seen[u] = -1;
 return res;
```

DominatorTree.h

Description: a dominates b iff every path from 1 to b passes through a 8cba41, 47 lines

```
template<int SZ> struct Dominator {
 vi adj[SZ], ans[SZ]; // input edges, edges of dominator tree
  vi radj[SZ], child[SZ], sdomChild[SZ];
  int label[SZ], rlabel[SZ], sdom[SZ], dom[SZ], co;
  int root = 1;
  int par[SZ], bes[SZ];
  int get(int x) {
   // DSU with path compression
    // get vertex with smallest sdom on path to root
   if (par[x] != x)
     int t = get(par[x]); par[x] = par[par[x]];
     if (sdom[t] < sdom[bes[x]]) bes[x] = t;</pre>
   return bes[x];
  void dfs(int x) { // create DFS tree
   label[x] = ++co; rlabel[co] = x;
    sdom[co] = par[co] = bes[co] = co;
   trav(y,adj[x]) {
     if (!label[y]) {
       dfs(y);
       child[label[x]].pb(label[y]);
     radj[label[y]].pb(label[x]);
  void init() {
   dfs(root);
   FORd(i,1,co+1) {
     trav(j,radj[i]) ckmin(sdom[i],sdom[get(j)]);
     if (i > 1) sdomChild[sdom[i]].pb(i);
     trav(j,sdomChild[i]) {
       int k = get(j);
       if (sdom[j] == sdom[k]) dom[j] = sdom[j];
       else dom[i] = k;
     trav(j,child[i]) par[j] = i;
   FOR(i, 2, co+1) {
     if (dom[i] != sdom[i]) dom[i] = dom[dom[i]];
     ans[rlabel[dom[i]]].pb(rlabel[i]);
};
```

EdgeColoring.h

Description: Vizing's Theorem: If max degree in simple undirected graph is d, then it can be edge colored with d+1 colors Time: ?

```
bd4b3d, 94 lines
template<int SZ> struct EdgeColor {
  int n, adjVert[SZ][SZ], adjCol[SZ][SZ];
  int deg[SZ], maxDeg;
  EdgeColor(int _n) {
   n = _n; maxDeq = 0;
   FOR(i,n) {
     deg[i] = 0;
     FOR(j,n) adjVert[i][j] = adjCol[i][j] = -1;
```

```
void delEdge(int x, int y) {
  if (adjVert[x][y] == -1) return;
  int C = adjVert[x][y];
  adjCol[x][C] = adjCol[y][C] = adjVert[x][y] = adjVert[y][x]
void setEdge(int x, int y, int c) { // delete previous value
  delEdge(x,y); assert(adjCol[x][c] == -1 && adjCol[y][c] ==
  adjVert[x][y] = adjVert[y][x] = c, adjCol[x][c] = y, adjCol
     \hookrightarrow[v][c] = x;
void shiftPath(int x, vi p) {
  FORd(i,sz(p)) setEdge(x,p[i],notAdj[p[i]]);
vi getPath(int st, int c0, int c1) {
  vi res = {st};
  for (int nex = 0; ; nex ^= 1) {
    int c = (nex == 0 ? c0 : c1);
    if (adjCol[res.back()][c] == -1) return res;
    res.pb(adjCol[res.back()][c]);
}
void flipPath(vi p, int c0, int c1) {
  FOR(i, sz(p)-1) delEdge(p[i],p[i+1]);
  FOR(i,sz(p)-1) {
    if (i&1) setEdge(p[i],p[i+1],c0);
    else setEdge(p[i],p[i+1],c1);
}
int notAdj[SZ];
void addEdge(int x, int y) {
  maxDeg = max(maxDeg, max(++deg[x], ++deg[y]));
  // generate a color which is not adjacent to each vertex
  FOR(i,n) {
    FOR(j, maxDeg+1) if (adjCol[i][j] == -1) {
      notAdj[i] = j;
      break;
  vi nex(n);
  FOR(i,n) if (adjVert[x][i] != -1) nex[i] = adjCol[x][notAdj]
     \hookrightarrow[i]];
  nex[y] = adjCol[x][notAdj[y]];
  // generate sequence of neighbors
  vi vis(n), seq = {y};
  while (seq.back() != -1 && !vis[seq.back()]) {
    vis[seq.back()] = 1;
    seq.pb(nex[seq.back()]);
  // case 1: easy
  if (seq.back() == -1) {
    seq.pop_back(), shiftPath(x,seq);
    return;
```

```
// separate into path and cycle
   int ind = 0; while (seq[ind] != seq.back()) ind ++;
   seq.pop_back();
   vi path = vi(seq.begin(), seq.begin()+ind);
   vi cyc = vi(seq.begin()+ind,seq.end());
   int c0 = notAdj[x], c1 = notAdj[cyc.back()];
   // case based on a/b path
   vi p = getPath(cyc.back(),c0,c1);
   if (p.back() != path.back()) {
     if (p.back() == x) { p.pop_back(), delEdge(x,p.back()); }
     flipPath(p,c0,c1);
     notAdj[seq.back()] = c0; shiftPath(x, seq);
    } else {
     reverse(all(p));
     flipPath(p,c0,c1);
     notAdj[path.back()] = c0; shiftPath(x,path);
};
```

Geometry (7)

7.1 Primitives

Point.h

```
Description: Easy Geo
                                                             708158, 47 lines
```

```
typedef ld T;
template \langle \text{class T} \rangle int \text{sqn}(\text{T x}) \{ \text{return } (\text{x} > 0) - (\text{x} < 0); \}
namespace Point {
 typedef pair<T,T> P;
 typedef vector<P> vP;
 P dir(T ang) {
    auto c = exp(ang*complex<T>(0,1));
    return P(c.real(),c.imag());
 T norm(P x) { return x.f*x.f*x.s*x.s; }
 T abs(P x) { return sqrt(norm(x)); }
 T angle(P x) { return atan2(x.s,x.f); }
 P conj(P x) \{ return P(x.f,-x.s); \}
 P operator+(const P& 1, const P& r) { return P(l.f+r.f,l.s+r.
  P operator-(const P& 1, const P& r) { return P(1.f-r.f,1.s-r.
  P operator* (const P& 1, const T& r) { return P(1.f*r,1.s*r);
 P operator*(const T& 1, const P& r) { return r*1; }
 P operator/(const P& 1, const T& r) { return P(1.f/r,1.s/r);
  P operator*(const P& 1, const P& r) { return P(1.f*r.f-1.s*r.
     \hookrightarrows,1.s*r.f+l.f*r.s); }
 P operator/(const P& 1, const P& r) { return 1*conj(r)/norm(r
     \hookrightarrow); }
  P& operator+=(P& 1, const P& r) { return 1 = 1+r; }
  P& operator = (P& 1, const P& r) { return 1 = 1-r; }
  P& operator*=(P& 1, const T& r) { return 1 = 1*r; }
 P& operator/=(P& 1, const T& r) { return 1 = 1/r; }
  P\& operator*=(P\& 1, const P\& r) { return 1 = 1*r; }
 P& operator/=(P& 1, const P& r) { return l = 1/r; }
 P unit(P x) { return x/abs(x); }
```

```
T dot(P a, P b) { return (conj(a)*b).f; }
T cross(P a, P b) { return (conj(a)*b).s; }
T cross(P p, P a, P b) { return cross(a-p,b-p); }
P rotate(P a, T b) { return a*P(cos(b),sin(b)); }

T dist(P p, P a, P b) { return std::abs(cross(p,a,b))/abs(a-b \( \to \)); }
P reflect(P p, P a, P b) { return a*conj((p-a)/(b-a))*(b-a); \( \to \)}
P foot(P p, P a, P b) { return (p*reflect(p,a,b))/(T)2; }
bool onSeg(P p, P a, P b) {
    return cross(a,b,p) == 0 && dot(p-a,p-b) <= 0;
}
};
using namespace Point;</pre>
```

AngleCmp.h

Description: sorts points according to atan2

fccaee, 5 lines

SegIntersect.h

Description: computes the intersection point(s) of line segments AB₁₅ CD₁₀

7.2 Polygons

Area h

Description: computes area + the center of mass of a polygon with constant mass per unit area $\frac{456403}{16}$ lines

```
T area(const vP& v) {
   T area = 0;
   FOR(i,sz(v)) {
    int j = (i+1)*sz(v); T a = cross(v[i],v[j]);
    area += a;
   }
   return std::abs(area)/2;
}
P centroid(const vP& v) {
   P cen(0,0); T area = 0; // 2*signed area
   FOR(i,sz(v)) {
    int j = (i+1)*sz(v); T a = cross(v[i],v[j]);
    cen += a*(v[i]+v[j]); area += a;
}
```

```
return cen/area/(T)3;
```

InPoly.h

Description: tests whether a point is inside, on, or outside the perimeter of any polygon 8f2d6a, 10 lines

```
string inPoly(const vP& p, P z) {
  int n = sz(p), ans = 0;
  FOR(i,n) {
    P x = p[i], y = p[(i+1)%n];
    if (onSeg(z,x,y)) return "on";
    if (x.s > y.s) swap(x,y);
    if (x.s <= z.s && y.s > z.s && cross(z,x,y) > 0) ans ^= 1;
  }
  return ans ? "in" : "out";
}
```

ConvexHull.h

 $\textbf{Description:} \ \, \textbf{Top-bottom convex hull}$

Time: $O(N \log N)$ 9be106, 37 lines

```
// typedef 11 T;
using namespace Point;
pair<vi, vi> ulHull(const vP& P) {
 vi p(sz(P)), u, l; iota(all(p), 0);
  sort(all(p), [&P](int a, int b) { return P[a] < P[b]; });</pre>
  trav(i,p) {
    #define ADDP(C, cmp) while (sz(C) > 1 \&\& cross(\
      P[C[sz(C)-2]], P[C.back()], P[i]) cmp 0) C.pop_back(); C.pb
         \hookrightarrow (i);
    ADDP (u, >=); ADDP (1, <=);
 return {u,1};
vi hullInd(const vP& P) {
 vi u, l; tie(u, l) = ulHull(P);
 if (sz(1) <= 1) return 1;
 if (P[1[0]] == P[1[1]]) return {0};
 1.insert (end(1), rbegin(u)+1, rend(u)-1); return 1;
vP hull(const vP& P) {
 vi v = hullInd(P);
 vP res; trav(t,v) res.pb(P[t]);
 return res;
ld diameter(vP P) { // rotating calipers
 P = hull(P);
 int n = sz(P), ind = 1; ld ans = 0;
 FOR(i,n)
    for (int j = (i+1) %n; ; ind = (ind+1) %n) {
      ckmax(ans, abs(P[i]-P[ind]));
      if (cross(P[j]-P[i],P[(ind+1)%n]-P[ind]) <= 0) break;</pre>
 return ans;
```

7.3 Circles

Circles.h

Description: misc operations with two circles

c4314f, 52 lines

```
using namespace Point;
```

```
namespace Circles {
 typedef pair<P,T> circ;
 bool on(circ x, P y) { return abs(y-x.f) == x.s; }
 bool in(circ x, P y) { return abs(y-x.f) <= x.s; }</pre>
 T arcLength(circ x, P a, P b) {
    P d = (a-x.f)/(b-x.f);
    return x.s*acos(d.f);
 P intersectPoint(circ x, circ y, int t = 0) { // assumes
     \hookrightarrow intersection points exist
    T d = abs(x.f-y.f); // distance between centers
    T theta = a\cos((x.s*x.s+d*d-y.s*y.s)/(2*x.s*d)); // law of
       \rightarrowcosines
    P tmp = (v.f-x.f)/d*x.s;
    return x.f+tmp*dir(t == 0 ? theta : -theta);
 T intersectArea(circ x, circ y) { // not thoroughly tested
    T d = abs(x.f-y.f), a = x.s, b = y.s; if (a < b) swap(a,b);
    if (d \ge a+b) return 0:
    if (d <= a-b) return PI*b*b;
    auto ca = (a*a+d*d-b*b)/(2*a*d), cb = (b*b+d*d-a*a)/(2*b*d)
    auto s = (a+b+d)/2, h = 2*sqrt(s*(s-a)*(s-b)*(s-d))/d;
    return a*a*acos(ca)+b*b*acos(cb)-d*h;
  P \text{ tangent}(P x, \text{ circ } y, \text{ int } t = 0)  {
    y.s = abs(y.s); // abs needed because internal calls y.s <
    if (v.s == 0) return v.f;
    T d = abs(x-v.f);
    P = pow(y.s/d, 2) * (x-y.f) + y.f;
    P b = sqrt(d*d-y.s*y.s)/d*y.s*unit(x-y.f)*dir(PI/2);
    return t == 0 ? a+b : a-b;
  vector<pair<P,P>> external(circ x, circ y) { // external
    \hookrightarrowtangents
    vector<pair<P,P>> v;
    if (x.s == y.s) {
      P tmp = unit(x.f-y.f)*x.s*dir(PI/2);
      v.pb(mp(x.f+tmp,y.f+tmp));
      v.pb(mp(x.f-tmp,y.f-tmp));
    } else {
      P p = (y.s*x.f-x.s*y.f)/(y.s-x.s);
      FOR(i,2) v.pb({tangent(p,x,i),tangent(p,y,i)});
    return v;
 vector<pair<P,P>> internal(circ x, circ y) { // internal
     \hookrightarrowtangents
    x.s \star = -1; return external (x, y);
using namespace Circles;
```

Circumcenter.h

Description: circumcenter

0d49ba, 5 lines

```
pair<P,T> ccCenter(P a, P b, P c) { // circumcenter, radius
  b -= a; c -= a;
  P res = b*c*(conj(c)-conj(b))/(b*conj(c)-conj(b)*c);
  return {a+res,abs(res)};
}
```

MinEnclosingCircle.h

Description: computes minimum enclosing circle

```
"Circumcenter.h" 63f976, 13 lines
pair<P, T> mec(vP ps) {
    shuffle(all(ps), mt19937(time(0)));
    P o = ps[0]; T r = 0, EPS = 1 + 1e-8;
    FOR(i,sz(ps)) if (abs(o-ps[i]) > r*EPS) {
        o = ps[i], r = 0;
        FOR(j,i) if (abs(o-ps[j]) > r*EPS) {
            o = (ps[i]+ps[j])/2, r = abs(o-ps[i]);
            FOR(k,j) if (abs(o-ps[k]) > r*EPS)
            tie(o,r) = ccCenter(ps[i],ps[j],ps[k]);
        }
    }
    return {o,r};
}
```

7.4 Misc

ClosestPair.h

Description: $O(N \log N)$ line sweep to find two closest points out of $N = N \log N$

```
using namespace Point;

pair<P,P> solve(vP v) {
    pair<ld,pair<P,P>> bes; bes.f = INF;
    set<P> S; int ind = 0;

sort(all(v));
    FOR(i,sz(v)) {
    if (i && v[i] == v[i-1]) return {v[i],v[i]};
    for (; v[i].f-v[ind].f >= bes.f; ++ind)
        S.erase({v[ind].s,v[ind].f});
    for (auto it = S.ub({v[i].s-bes.f,INF});
        it != end(S) && it->f < v[i].s+bes.f; ++it) {
        P t = {it->s,it->f};
        ckmin(bes,{abs(t-v[i]),{t,v[i]}});
    }
    S.insert({v[i].s,v[i].f});
}
return bes.s;
}
```

DelaunavFast.h

 $\mathbf{Description:}$ Delaunay Triangulation, concyclic points are OK (but not all collinear)

```
"Point.h"
                                                      765ba9, 94 lines
typedef 11 T:
typedef struct Quad* Q;
typedef __int128_t 111; // (can be 11 if coords are < 2e4)
P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
struct Quad {
  bool mark; Q o, rot; P p;
  P F() { return r()->p; }
  Q r() { return rot->rot; }
  Q prev() { return rot->o->rot;
  Q next() { return r()->prev(); }
// test if p is in the circumcircle
bool circ(P p, P a, P b, P c) {
  ll ar = cross(a,b,c); assert(ar); if (ar < 0) swap(a,b);
  111 p2 = norm(p), A = norm(a) - p2,
    B = norm(b) - p2, C = norm(c) - p2;
```

```
return cross(p,a,b) *C+cross(p,b,c) *A+cross(p,c,a) *B > 0;
Q makeEdge(P orig, P dest) {
  Q q[] = \{new Quad\{0,0,0,orig\}, new Quad\{0,0,0,arb\},
       new Quad{0,0,0,dest}, new Quad{0,0,0,arb}};
  FOR(i, 4) q[i] -> o = q[-i \& 3], q[i] -> rot = q[(i+1) \& 3];
  return *q;
void splice(O a, O b) {
  swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
0 connect(0 a, 0 b) {
  Q = makeEdge(a->F(), b->p);
  splice(q, a->next());
  splice(q->r(), b);
  return q;
pair<Q,Q> rec(const vector<P>& s) {
  if (sz(s) \le 3) {
    Q = makeEdge(s[0], s[1]), b = makeEdge(s[1], s.back());
    if (sz(s) == 2) return { a, a->r() };
    splice(a->r(), b);
    auto side = cross(s[0], s[1], s[2]);
    0 c = side ? connect(b, a) : 0;
    return {side < 0 ? c->r() : a, side < 0 ? c : b->r() };
#define H(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (cross(e->F(),H(base)) > 0)
 O A, B, ra, rb;
 int half = sz(s) / 2;
 tie(ra, A) = rec({all(s) - half});
  tie(B, rb) = rec({sz(s) - half + all(s)});
  while ((cross(B->p,H(A)) < 0 && (A = A->next()))
       (cross(A->p,H(B)) > 0 && (B = B->r()->o)));
  O base = connect(B->r(), A);
  if (A->p == ra->p) ra = base->r();
  if (B->p == rb->p) rb = base;
#define DEL(e, init, dir) Q e = init->dir; if (valid(e)) \
    while (circ(e->dir->F(), H(base), e->F())) { \
      Q t = e->dir; \setminus
      splice(e, e->prev()); \
      splice(e->r(), e->r()->prev()); \
      e = t; \
  for (;;) {
    DEL(LC, base->r(), o); DEL(RC, base, prev());
    if (!valid(LC) && !valid(RC)) break;
    if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
      base = connect(RC, base->r());
      base = connect(base->r(), LC->r());
  return {ra, rb};
vector<array<P,3>> triangulate(vector<P> pts) {
  sort(all(pts)); assert(unique(all(pts)) == pts.end());
  if (sz(pts) < 2) return {};
  Q = rec(pts).f; vector < Q > q = {e};
  int qi = 0;
  while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o;
#define ADD { Q c = e; do { c->mark = 1; pts.push_back(c->p); \
  q.push\_back(c->r()); c = c->next(); } while (c != e); }
  ADD; pts.clear();
```

```
while (qi < sz(q)) if (!(e = q[qi++])->mark) ADD;
vector<array<P,3>> ret;
FOR(i,sz(pts)/3) ret.pb({pts[3*i],pts[3*i+1],pts[3*i+2]});
return ret;
}
```

$7.5 \quad 3D$

Point3D.h

```
Description: Basic 3D Geometry
```

a4d471, 45 lines

```
typedef ld T;
namespace Point3D {
 typedef array<T,3> P3;
 typedef vector<P3> vP3:
 T norm(const P3& x) {
   T sum = 0; FOR(i,sz(x)) sum += x[i]*x[i];
   return sum;
 T abs(const P3& x) { return sqrt(norm(x)); }
 P3& operator+=(P3& 1, const P3& r) { F0R(i,3) 1[i] += r[i];
     →return 1; }
 P3& operator = (P3& 1, const P3& r) { F0R(i,3) 1[i] -= r[i];
     →return 1; }
 P3& operator*=(P3& 1, const T& r) { F0R(i,3) 1[i] *= r;
     →return 1; }
 P3& operator/=(P3& 1, const T& r) { F0R(i,3) 1[i] /= r;
     →return 1: }
 P3 operator+(P3 1, const P3& r) { return 1 += r; }
 P3 operator-(P3 1, const P3& r) { return 1 -= r; }
 P3 operator*(P3 1, const T& r) { return 1 *= r; }
 P3 operator*(const T& r, const P3& 1) { return 1*r; }
 P3 operator/(P3 1, const T& r) { return 1 /= r; }
 T dot(const P3& a, const P3& b) {
   T sum = 0; FOR(i,3) sum += a[i]*b[i];
   return sum;
 P3 cross(const P3& a, const P3& b) {
   return {a[1]*b[2]-a[2]*b[1],
       a[2]*b[0]-a[0]*b[2],
       a[0]*b[1]-a[1]*b[0];
 bool isMult(const P3& a, const P3& b) {
   auto c = cross(a,b);
   FOR(i,sz(c)) if (c[i] != 0) return 0;
   return 1;
 bool collinear (const P3& a, const P3& b, const P3& c) {
     →return isMult(b-a,c-a); }
 bool coplanar(const P3& a, const P3& b, const P3& c, const P3
    S d) {
    return isMult(cross(b-a,c-a),cross(b-a,d-a));
using namespace Point3D;
```

Hull3D.h

Description: 3D Convex Hull + Polyedron Volume **Time:** $\mathcal{O}(N^2)$

d201e5, 50 lines

2337c9, 36 lines

```
using namespace Point3D;
struct ED {
  void ins(int x) { (a == -1 ? a : b) = x; }
  void rem(int x) { (a == x ? a : b) = -1; }
  int cnt() { return (a !=-1) + (b !=-1); }
  int a, b;
struct F { P3 q; int a, b, c; };
vector<F> hull3d(const vP3& A) {
  assert(sz(A) >= 4);
  vector<vector<ED>> E(sz(A), vector<ED>(sz(A), \{-1, -1\}));
  #define E(x,y) E[f.x][f.y]
  vector<F> FS; // faces
  auto mf = [&] (int i, int j, int k, int l) { // make face
    P3 q = cross(A[\dot{j}]-A[\dot{i}], A[\dot{k}]-A[\dot{i}]);
    if (dot(q, A[1]) > dot(q, A[i])) q *= -1; // make sure q
       \hookrightarrowpoints outward
    F f{q, i, j, k};
    E(a,b).ins(k); E(a,c).ins(j); E(b,c).ins(i);
    FS.pb(f);
  FOR(i, 4) FOR(j, i+1, 4) FOR(k, j+1, 4) mf(i, j, k, 6-i-j-k);
  FOR(i, 4, sz(A)) {
    FOR(j,sz(FS)) {
      F f = FS[j];
      if (dot(f,q,A[i]) > dot(f,q,A[f,a]))  { // face is visible
         \hookrightarrow, remove edges
        E(a,b).rem(f.c), E(a,c).rem(f.b), E(b,c).rem(f.a);
        swap(FS[j--], FS.back());
        FS.pop_back();
    FOR(j,sz(FS)) { // add faces with new point
      F f = FS[j];
      \#define C(a, b, c) if (E(a,b).cnt() != 2) mf(f.a, f.b, i, b)
         \hookrightarrow f.c);
      C(a, b, c); C(a, c, b); C(b, c, a);
  trav(it, FS) if (dot(cross(A[it.b]-A[it.a], A[it.c]-A[it.a]),
     \hookrightarrowit.q) <= 0)
    swap(it.c, it.b);
  return FS;
} // computes hull where no four are coplanar
T signedPolyVolume(const vP3& p, const vector<F>& trilist) {
 T v = 0:
  trav(i,trilist) v += dot(cross(p[i.a],p[i.b]),p[i.c]);
  return v/6;
```

Strings (8)

8.1 Lightweight

KMP.h

Description: f[i] equals the length of the longest proper suffix of the i-th prefix of s that is a prefix of s

```
Time: \mathcal{O}\left(N\right) 08f252, 15 lines
```

```
vi kmp(string s) {
  int N = sz(s); vi f(N+1); f[0] = -1;
```

```
FOR(i,1,N+1) {
    f[i] = f[i-1];
    while (f[i] != -1 && s[f[i]] != s[i-1]) f[i] = f[f[i]];
    f[i] ++;
}
    return f;
}

vi getOc(string a, string b) { // find occurrences of a in b
    vi f = kmp(a+"@"+b), ret;
    FOR(i,sz(a),sz(b)+1) if (f[i+sz(a)+1] == sz(a)) ret.pb(i-sz(a \( \to ) \));
    return ret;
}
```

Z.h

Description: similar to KMP **Time:** O(N)

```
a4e01c, 19 lines
vi z(string s) {
 int N = sz(s); s += '#';
 vi ans(N); ans[0] = N;
 int L = 1, R = 0;
 FOR(i,1,N) {
   if (i \le R) ans[i] = min(R-i+1, ans[i-L]);
   while (s[i+ans[i]] == s[ans[i]]) ans[i] ++;
   if (i+ans[i]-1 > R) L = i, R = i+ans[i]-1;
 return ans;
vi getPrefix(string a, string b) { // find prefixes of a in b
 vi t = z(a+b), T(sz(b));
 FOR(i,sz(T)) T[i] = min(t[i+sz(a)],sz(a));
 return T:
// pr(z("abcababcabcaba"), getPrefix("abcab", "uwetrabcerabcab"))
  \hookrightarrow;
```

Manacher.h

Description: Calculates length of largest palindrome centered at each character of string **Time:** $\mathcal{O}(N)$

```
34a78b, 18 lines
vi manacher(string s) {
 string s1 = "@";
 trav(c,s) s1 += c, s1 += "#";
 s1[sz(s1)-1] = '&';
 vi ans (sz(s1)-1);
 int 10 = 0, hi = 0;
 FOR(i, 1, sz(s1) - 1) {
    if (i != 1) ans[i] = min(hi-i, ans[hi-i+lo]);
    while (s1[i-ans[i]-1] == s1[i+ans[i]+1]) ans[i] ++;
    if (i+ans[i] > hi) lo = i-ans[i], hi = i+ans[i];
 ans.erase(begin(ans));
 FOR(i,sz(ans)) if ((i&1) == (ans[i]&1)) ans[i] ++; // adjust
    \hookrightarrow lengths
 return ans;
// ps(manacher("abacaba"))
```

MinRotation.h

Description: minimum rotation of string **Time:** $\mathcal{O}(N)$

```
483a1a, 8 lines
```

LyndonFactorization.h

Description: A string is "simple" if it is strictly smaller than any of its own nontrivial suffixes. The Lyndon factorization of the string s is a factorization $s = w_1 w_2 \dots w_k$ where all strings w_i are simple and $w_1 \ge w_2 \ge \dots \ge w_k$ **Time:** $\mathcal{O}(N)$

```
vector<string> duval(const string& s) {
 int n = sz(s); vector<string> factors;
 for (int i = 0; i < n; ) {
    int j = i + 1, k = i;
    for (; j < n \&\& s[k] <= s[j]; j++) {
     if (s[k] < s[j]) k = i;
      else k ++;
    for (; i \le k; i += j-k) factors.pb(s.substr(i, j-k));
 return factors;
int minRotation(string s) { // get min index i such that cyclic
  \hookrightarrow shift starting at i is min rotation
 int n = sz(s); s += s;
 auto d = duval(s); int ind = 0, ans = 0;
  while (ans+sz(d[ind]) < n) ans += sz(d[ind++]);
  while (ind && d[ind] == d[ind-1]) ans -= sz(d[ind-1]);
 return ans;
```

8.2 Suffix Structures

ACfixed.h

Description: for each prefix, stores link to max length suffix which is also a prefix

```
Time: \mathcal{O}(N \Sigma)
struct ACfixed { // fixed alphabet
 struct node {
    array<int,26> to;
   int link;
  };
 vector<node> d;
 ACfixed() { d.emplace_back(); }
 int add(string s) { // add word
    int v = 0;
    trav(C,s)
      int c = C-'a';
      if (!d[v].to[c]) {
        d[v].to[c] = sz(d);
        d.emplace_back();
      v = d[v].to[c];
    return v;
```

```
void init() { // generate links
   d[0].link = -1;
   queue<int> q; q.push(0);
   while (sz(q)) {
      int v = q.front(); q.pop();
      FOR(c,26) {
       int u = d[v].to[c]; if (!u) continue;
        d[u].link = d[v].link == -1 ? 0 : d[d[v].link].to[c];
        q.push(u);
      }
      if (v) FOR(c,26) if (!d[v].to[c])
        d[v].to[c] = d[d[v].link].to[c];
    }
};
```

PalTree.h

 $\textbf{Description:} \ \ \mathrm{palindromic} \ \ \mathrm{tree}$

 $\mathbf{Time:}\ \mathcal{O}\left(N\right)$

36a5a4, 26 lines

```
template<int SZ> struct PalTree {
  static const int sigma = 26;
  int s[SZ], len[SZ], link[SZ], to[SZ][sigma], oc[SZ];
  int n, last, sz;
  PalTree() { s[n++] = -1; link[0] = 1; len[1] = -1; sz = 2; }
  int getLink(int v) {
   while (s[n-len[v]-2] != s[n-1]) v = link[v];
    return v:
  void addChar(int c) {
   s[n++] = c;
   last = getLink(last);
   if (!to[last][c]) {
     len[sz] = len[last]+2;
     link[sz] = to[getLink(link[last])][c];
     to[last][c] = sz++;
    last = to[last][c]; oc[last] ++;
  void init() { // number of occurrences of each palindrome
   vpi v; FOR(i,2,sz) v.pb({len[i],i});
   sort(all(v)); reverse(all(v));
   trav(a,v) oc[link[a.s]] += oc[a.s];
};
```

SuffixArray.h Description: ?

Description: ? **Time:** $O(N \log N)$

dbc6b9, 51 lines

```
template<int SZ> struct SuffixArray {
  string S; int N;
 void init(const string& _S) {
   S = _S; N = sz(S);
   genSa(); genLcp();
    // R.init(lcp);
  vi sa, isa;
  void genSa() { // http://ekzlib.herokuapp.com
    sa.rsz(N); vi classes(N);
   FOR(i,N) sa[i] = N-1-i, classes[i] = S[i];
   stable_sort(all(sa), [this](int i, int j) { return S[i] < S</pre>
      \hookrightarrow[j]; });
    for (int len = 1; len < N; len *= 2) {
     vi c(classes);
     FOR(i,N) { // compare first len characters of each suffix
       bool same = i \&\& sa[i-1] + len < N
```

```
&& c[sa[i]] == c[sa[i-1]]
                && c[sa[i]+len/2] == c[sa[i-1]+len/2];
        classes[sa[i]] = same ? classes[sa[i-1]] : i;
      vi nex(N), s(sa); iota(all(nex),0); // suffixes with <=
         →len chars will not change pos
      FOR(i,N) {
       int s1 = s[i]-len;
       if (s1 \ge 0) sa[nex[classes[s1]]++] = s1; // order
           →pairs w/ same first len chars by next len chars
    isa.rsz(N); FOR(i,N) isa[sa[i]] = i;
 vi lcp;
 void genLcp() { // KACTL
   lcp = vi(N-1);
    int h = 0;
   FOR(i,N) if (isa[i]) {
     int pre = sa[isa[i]-1];
      while (\max(i,pre)+h < N \&\& S[i+h] == S[pre+h]) h++;
      lcp[isa[i]-1] = h; // lcp of suffixes starting at pre and
      if (h) h--: // if we cut off first chars of two strings
        \hookrightarrowwith lcp h, then remaining portions still have lcp h
 }
 /*RMO<int,SZ> R;
 int getLCP(int a, int b) {
   if (max(a,b) >= N) return 0;
   if (a == b) return N-a;
   int t0 = isa[a], t1 = isa[b];
   if (t0 > t1) swap(t0,t1);
   return R.query(t0,t1-1);
 } */
};
```

ReverseBW.h

Description: Reverse Burrows-Wheeler

13b6b0, 8 lines

```
string reverseBW(string s) {
  int nex[sz(s)];
  vector<pair<char,int>> v; FOR(i,sz(s)) v.pb({s[i],i});
  sort(all(v)); FOR(i,sz(v)) nex[i] = v[i].s;
  int cur = nex[0]; string ret;
  for (; cur;cur = nex[cur]) ret += v[cur].f;
  return ret;
}
```

SuffixAutomaton.h

Description: Suffix Automaton **Time:** $\mathcal{O}(N \log \Sigma)$

struct SuffixAutomaton {
 struct state {
 int len = 0, firstPos = -1, link = -1;
 bool isClone = 0;
 map<char, int> next;
 vi invLink;
 };

 vector<state> st;
 int last = 0;
 void extend(char c) {
 int cur = s2(st); st.eb();
}

```
st[cur].len = st[last].len+1, st[cur].firstPos = st[cur].
       \hookrightarrowlen-1:
    int p = last;
    while (p != -1 && !st[p].next.count(c)) {
      st[p].next[c] = cur;
      p = st[p].link;
    if (p == -1) {
      st[cur].link = 0;
    } else {
      int q = st[p].next[c];
      if (st[p].len+1 == st[q].len) {
       st[cur].link = q;
      } else {
       int clone = sz(st); st.pb(st[q]);
        st[clone].len = st[p].len+1, st[clone].isClone = 1;
        while (p != -1 \&\& st[p].next[c] == q) {
          st[p].next[c] = clone;
          p = st[p].link;
        st[q].link = st[cur].link = clone;
   last = cur;
 void init(string s) {
   st.eb(); trav(x,s) extend(x);
   FOR(v,1,sz(st)) st[st[v].link].invLink.pb(v);
 void getAllOccur(vi& oc, int v) {
   if (!st[v].isClone) oc.pb(st[v].firstPos);
   trav(u,st[v].invLink) getAllOccur(oc,u);
 vi allOccur(string s) {
   int cur = 0;
   trav(x,s) {
     if (!st[cur].next.count(x)) return {};
      cur = st[cur].next[x];
   vi oc; qetAllOccur(oc, cur); trav(t, oc) t += 1-sz(s);
   sort(all(oc)); return oc;
 vl distinct;
 11 getDistinct(int x) {
   if (distinct[x]) return distinct[x];
    distinct[x] = 1;
   trav(y,st[x].next) distinct[x] += getDistinct(y.s);
   return distinct[x];
 ll numDistinct() { // # of distinct substrings, including
     \hookrightarrowempty
    distinct.rsz(sz(st));
   return getDistinct(0);
 ll numDistinct2() { // another way to get # of distinct
     \hookrightarrow substrings
    11 \text{ ans} = 1;
   FOR(i, 1, sz(st)) ans += st[i].len-st[st[i].link].len;
   return ans;
};
```

SuffixTree.h Description: ? Time: $\mathcal{O}(N \Sigma)$

61394a, 50 lines

```
struct SuffixTree {
```

```
enum { N = 200010, ALPHA = 26 }; // N \sim 2*maxlen+10
 int toi(char c) { return c - 'a'; }
  string a; // v = cur node, g = cur position
 int t[N][ALPHA], 1[N], r[N], p[N], s[N], v=0, q=0, m=2;
  void ukkadd(int i, int c) { suff:
   if (r[v] <=q) {
     if (t[v][c]==-1) { t[v][c]=m; l[m]=i;
       p[m++]=v; v=s[v]; q=r[v]; qoto suff; }
     v=t[v][c]; q=l[v];
    if (q==-1 || c==toi(a[q])) q++; else {
     l[m+1]=i; p[m+1]=m; l[m]=l[v]; r[m]=q;
     p[m]=p[v]; t[m][c]=m+1; t[m][toi(a[q])]=v;
     l[v]=q; p[v]=m; t[p[m]][toi(a[l[m]])]=m;
     v=s[p[m]]; q=l[m];
     while (q<r[m]) { v=t[v][toi(a[q])]; q+=r[v]-l[v]; }</pre>
     if (q==r[m]) s[m]=v; else s[m]=m+2;
     q=r[v]-(q-r[m]); m+=2; qoto suff;
  SuffixTree(string a) : a(a) {
   fill(r,r+N,sz(a));
   memset(s, 0, sizeof s);
   memset(t, -1, sizeof t);
   fill(t[1],t[1]+ALPHA,0);
   s[0] = 1; 1[0] = 1[1] = -1; r[0] = r[1] = p[0] = p[1] = 0;
   FOR(i,sz(a)) ukkadd(i, toi(a[i]));
  // example: find longest common substring (uses ALPHA = 28)
 pi best;
  int lcs(int node, int i1, int i2, int olen) {
   if (1[node] <= i1 && i1 < r[node]) return 1;</pre>
   if (1[node] <= i2 && i2 < r[node]) return 2;</pre>
   int mask = 0, len = node ? olen + (r[node] - l[node]) : 0;
   FOR(c, ALPHA) if (t[node][c] != -1)
     mask |= lcs(t[node][c], i1, i2, len);
   if (mask == 3)
     best = max(best, {len, r[node] - len});
   return mask;
  static pi LCS(string s, string t) {
   SuffixTree st(s + (char) ('z' + 1) + t + (char) ('z' + 2));
   st.lcs(0, sz(s), sz(s) + 1 + sz(t), 0);
   return st.best;
};
```

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