# **Binary Interpolation Tree (B.I.T)**

Binary Interpolation Tree is a machine learning model i have been working on mostly out of curiosity. Please let me know if my terminology is off, i made any other errors, or something very similar already exists.

The main idea was to conceptionalize a model that provides:

- fast runtime on cpu's in relation to model complexity (O(n),  $\Omega(1)$ , input dependent)
- the ability to further simplify computations based on a maximum inaccuracy
- simple and fast training
- very low initial model complexity that increases during training
- multithreading capabilities for maximum cpu usage
- increased transparency compared to other ML models
- a parameter regarding determinism of the algorithm

### 1. The Model

#### 1.1 The basic model

As the name suggests, this model is basically a binary tree. This binary tree consists of nodes and end-nodes. Nodes always have 2 distinct children except for end-nodes wich have non. Every node has a reference to its parent, except for the root-node wich has no parent. Every node has a position vector and every end-node additionally contains an output vector. For the purpose of describing this system mathematically, every node has an index i and the following variables associated with it:

The index of the first child:  $c_1(i)$ .

The index of the second child:  $c_2(i)$ .

The index of the parent: parent(i).

The position vector:  $\overrightarrow{p}(i) \in V_i$ .

The result of the calculation for nodes, or the output vector for end-nodes:  $\overrightarrow{o}(i) \in V_o$ .

For the root-node the index i = 0.

An index of -1 means there is no reference. For example parent(0) := -1.

To calculate the output  $\overrightarrow{o}(0)$  of this model for a given input vector  $\overrightarrow{x} \in V_i$ , this binary tree is traversed recursively in my cpu implementation, from the root-node, to the end-nodes. The following calculations are performed for every node:

$$I(\overrightarrow{p_1},\overrightarrow{p_2},\overrightarrow{x}):=\frac{(\overrightarrow{p_2}-\overrightarrow{p_1})\circ(\overrightarrow{x}-\overrightarrow{p_1})}{\left|\overrightarrow{p_2}-\overrightarrow{p_1}\right|^2} \qquad clamp(x):=\min(1,\max(0,x))$$

$$a(c_1(i)) := \operatorname{clamp}(I(\overrightarrow{p}(c_1(i)), \overrightarrow{p}(c_2(i)), \overrightarrow{x})) \qquad \qquad a(c_2(i)) := 1 - a(c_1(i))$$

$$\overrightarrow{o}(i) := a(c_1(i)) \cdot \overrightarrow{o}(c_1(i)) + a(c_2(i)) \cdot \overrightarrow{o}(c_2(i))$$

if  $a(c_1(i))=0$  or  $a(c_2(i))=0$ ,  $\overrightarrow{o}(c_1(i))$  or  $\overrightarrow{o}(c_2(i))$  doesn't have to be calculated, wich can cause a great performance increase for sequential execution. the following statements are true and fully define  $I(\overrightarrow{p_1},\overrightarrow{p_2},\overrightarrow{x})$  in an intuitive way:

$$I(\overrightarrow{p_1}, \overrightarrow{p_2}, a * \overrightarrow{p_1} + (1 - a) * \overrightarrow{p_2}) = 1 - a \Rightarrow I(\overrightarrow{p_1}, \overrightarrow{p_2}, \overrightarrow{p_1}) = 0 \land I(\overrightarrow{p_1}, \overrightarrow{p_2}, \overrightarrow{p_2}) = 1$$
  
$$I(\overrightarrow{p_1}, \overrightarrow{p_2}, \overrightarrow{x_1}) = I(\overrightarrow{p_1}, \overrightarrow{p_2}, \overrightarrow{x_2}) \Leftrightarrow (\overrightarrow{p_2} - \overrightarrow{p_1}) \circ (\overrightarrow{x_1} - \overrightarrow{x_2}) = 0.$$

#### 1.2 Introducing indeterministic behavior

Swapping the definition of  $a(c_1(i))$  from  $a(c_1(i)) := clamp(I(\overrightarrow{p}(c_1(i)), \overrightarrow{p}(c_2(i)), \overrightarrow{x}))$  to  $a(c_1(i)) := clamp(I(\overrightarrow{p}(c_1(i)), \overrightarrow{p}(c_2(i)), \overrightarrow{x}) * d + (1 - d) * r)$ 

where  $0 \le d \le 1 \land 0 \le r \le 1$ , will allow indeterministic behavior, if a random number is assigned to r for each computation of  $a(c_1(i))$ . The parameter d is called determinism in my implementation, because d=1 means that the model is 100% deterministic while d=0 means the interpolation factor  $a(c_1(i))$  is completely random between 0 and 1.

#### 1.2 Simplifying computations by defining a maximum inaccuracy

For this purpose, the following variables are defined:

The number of dimensions of the input space:  $m := \dim(V_i)$ .

The number of dimensions of the output space:  $n := \dim(V_o)$ .

The Inaccuracy-scale vector  $\vec{s} \in V_o$  allows different maximum inaccuracies for all different output dimensions and is meant to be immutable for every model. The maximum inaccuracy t is meant to be passed for every Forward pass, to be able to calculate the output for the same model with varying speed and accuracy. The following variables are defined for each node:

The complexity:  $complexity(i) := 1 + complexity(c_1(i)) + complexity(c_2(i))$  for all nodes and complexity(i) := 1 for all end-nodes specifically.

complexity(i) is roughly proportional to the runtime of the forward pass off the node with the index i.

The contribution to the final result: 
$$c(i) := \begin{cases} c(parent(i)) * a(i) & i \neq 0 \\ 1 & i = 0 \end{cases}$$

The maximum contribution value 
$$v(i) := \max_{1 \leq j \leq n} (\frac{\overrightarrow{o}(i)_j}{\overrightarrow{s}_j})$$
 for all end-nodes and

 $v(i) := \max(v(c_1(i)), v(c_2(i)))$  for every node that is no end-node. v(i) can be cached for every node, because  $\vec{s}$  is meant to be immutable for every model.

$$\sum_{k \in M} v(k) \cdot c(k) \le t \Rightarrow \forall 1 \le j \le n : \sum_{k \in M} \overrightarrow{o}(k)_j * c(k) \le \overrightarrow{s}_j * t$$

This Expression means, that you can set  $\overrightarrow{o}(k) := 0$  for  $k \in M$  and can guarantee that the inaccuracy of  $\overrightarrow{o}(0)$ , caused by this simplification, will be less than  $\overrightarrow{s}*t$  for each output dimension, wich is a way of trading accuracy for speed. In my c# implementation, i wrote the recursive forward pass as an iterator method, to collect all nodes where  $v(i) \cdot c(i) \leq t$  (potential minors) and choose as much nodes as possible,

that have the lowest  $\frac{v(i) \cdot c(i)}{complexity(i)}$  ratio from those potential minors to be in M, so that

$$\sum_{k \in M} v(k) \cdot c(k) \le t$$
 is still satisfied. All nodes in  $M$  are called minors and get selected by

$$\frac{v(i) \cdot c(i)}{complexity(i)}$$
 because of the following statements:

 $v(i_1) \cdot c(i_1) < v(i_2) \cdot c(i_2) \wedge complexity(i_1) = complexity(i_2)$  means that it is beneficial to choose the node with the index  $i_1$  as minor, because this introduces less in inaccuracy, but roughly the same runtime.

 $v(i_1) \cdot c(i_1) = v(i_2) \cdot c(i_2) \wedge complexity(i_1) < complexity(i_2)$  means that it is beneficial to choose the node with the index  $i_1$  as minor, because this introduces the same inaccuracy, but it tends run the forward pass faster.

 $v(i_1) \cdot c(i_1) + v(i_2) \cdot c(i_2) = v(i_3) \cdot c(i_3) \wedge complexity(i_1) + complexity(i_2) = complexity(i_3)$  means that choosing the nodes with the index  $i_1$  and  $i_2$  as minor, will have almost the same impact as choosing only the node with the index  $i_3$ , because this introduces the same inaccuracy, and roughly the same runtime.

## 2. Supervised training of the model

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