Peter Smith, Introduction to Formal Logic (CUP, 2nd edition)

Exercises 33: More QL proofs

Do the unstarred examples in both (a) and (b) and check your answers before returning to the starred ones.

- (a) As a warm-up exercise, consider which of the following QL arguments ought to be valid (assume the wffs are interpreted). Give proofs warranting the valid inferences.
 - (1) $\exists x Sxxx : \exists x \exists y \exists z Sxyz$
 - (2) $\exists x \forall y \forall y Sxyy : \exists x Sxxx$
 - (3) $\forall x \exists y \forall z Sxyz : \exists x \forall y \exists z Sxzy$
- (4^*) $\neg \exists x \forall y \forall z Sxyz :: \forall x \exists z \exists y \neg Sxyz$
- $(5^*) \neg \exists x (\mathsf{Fx} \land \exists y (\mathsf{Gy} \land \mathsf{Lxy})) \therefore \forall x \forall y (\mathsf{Fx} \rightarrow (\mathsf{Gy} \rightarrow \neg \mathsf{Lxy}))$
- (b) Render the following inferences into suitable QL languages and provide derivations of the conclusions from the premisses in each case:
 - (1) Some people are boastful. No one likes anyone boastful. Therefore some people aren't liked by anyone.
 - (2) There's someone such that if *they* admire some philosopher, then I'm a Dutchman. So if *everyone* admires some philosopher, then I'm a Dutchman.
 - (3) Some good philosophers admire Frank; all wise people admire any good philosopher; Frank is wise; hence there is someone who both admires and is admired by Frank.
 - (4) Everyone loves themself if there's someone who loves them or whom they love. There's someone who is loved. Therefore someone loves themself.
- (5*) Only rational people with good judgement are logicians. Those who take some creationist myth literally lack good judgement. So logicians do not take any creationist myth literally.
- (6*) Given any two people, if the first admires Gödel and Gödel admires the second, then the first admires the second. Gödel admires anyone who has understood *Principia*. There's someone who has understood *Principia* who admires Gödel. Therefore there's someone who has understood *Principia* who admires everyone who has understood *Principia*!
- (7*) Any adult elephant weighs more than any horse. Some horse weighs more than any donkey. If a first thing weighs more than a second thing, and the second thing weighs more than a third, then the first weighs more than the third. Hence any adult elephant weighs more than any donkey.
- (c) Why should the following QL wffs be logically true (assuming the wffs are interpreted)?
 - (1) $\exists x(Fx \rightarrow \forall yFy)$
 - (2) $\forall x \exists y (\exists z Lxz \rightarrow Lxy)$
 - $(3) \quad \exists \mathsf{x} \forall \mathsf{y} (\neg \mathsf{F} \mathsf{y} \vee \mathsf{F} \mathsf{x})$

Show that those wffs are QL theorems – feel free now to use the derived rules $(\neg \forall)$ and $(\neg \exists)$ from §32.5 and to skip PL reasoning. Also give derivations to warrant the following inferences:

- (4) $(\forall x Fx \rightarrow \exists y Gy) : \exists x \exists y (Fx \rightarrow Gy)$
- (5) $(\exists z Fz \rightarrow \exists z Gz) \therefore \forall x \exists y (Fx \rightarrow Gy)$
- (6) $\forall x \exists y (Fy \rightarrow Gx) : \exists y \forall x (Fy \rightarrow Gx)$
- (d*) Although all our examples of QL proofs so far start from zero or more sentences and end with a sentence, we won't build that into our official characterization of QL proofs they can go from wffs involving dummy names to a conclusion which may involve a dummy name.

Say that the wffs Γ (not necessarily all sentences) are QL-consistent if there is no QL proof using wffs Γ as premisses and ending with ' \bot '; otherwise Γ are QL-inconsistent – compare Exercises 22(d*).

Assuming that the terms mentioned belong to the relevant language, show

(1) If the wffs Γ , $\alpha(\tau)$ are QL-inconsistent and the wffs Γ include $\forall \xi \alpha(\xi)$ then those wffs Γ are already QL-inconsistent.

and then conclude that

(2) If the wffs Γ are QL-consistent and $\forall \xi \alpha(\xi)$ is one of them, then $\Gamma, \alpha(\tau)$ are also QL-consistent.

Show further that

(2*) If the wffs Γ are QL-consistent and $\forall \xi \alpha(\xi)$ is one of them, then $\Gamma, \alpha(\tau_1), \alpha(\tau_2), \ldots, \alpha(\tau_k)$ all together are also QL-consistent (for any terms $\tau_1, \tau_2, \ldots, \tau_k$ of the relevant language).

Show similarly that

(3) If the wffs Γ are QL-consistent and $\exists \xi \alpha(\xi)$ is one of them, then $\Gamma, \alpha(\delta)$ are also QL-consistent if δ is a dummy name that doesn't appear in Γ .

Also show that

- (4) If the wffs Γ are QL-consistent and $\neg \forall \xi \alpha(\xi)$ is one of them, then $\Gamma, \exists \xi \neg \alpha(\xi)$ are QL-consistent.
- (5) If the wffs Γ are QL-consistent and $\neg \exists \xi \alpha(\xi)$ is one of them, then $\Gamma, \forall \xi \neg \alpha(\xi)$ are QL-consistent.