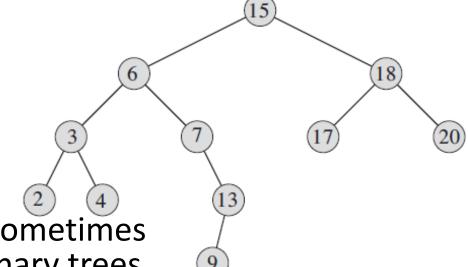
Binary Search Tree

Introduction

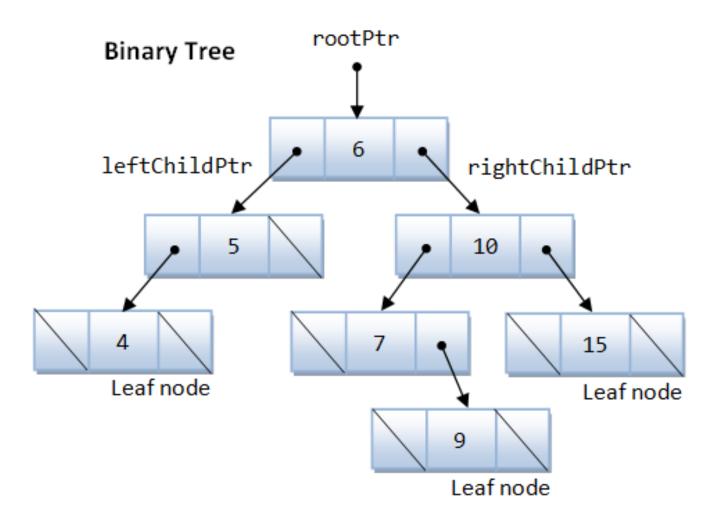


- Binary search trees (BST), sometimes called ordered or sorted binary trees.
- Properties:
 - The left subtree of a node contains only nodes with keys less than the node's key.
 - The right subtree of a node contains only nodes with keys greater than the node's key.
 - The left and right subtree each must also be a binary search tree. There must be no duplicate nodes.
- BSTs keep keys in sorted order, so that lookup, addition, and deletion operations can be executed efficiently.

Contd...

- All the operations traverse the tree from root to leaf, making comparisons to keys stored in the nodes of the tree and deciding, based on the comparison, to continue searching in the left or right subtrees.
- On average, this means that each comparison allows the operations to skip about half of the tree, so that each lookup, insertion or deletion takes time proportional to the logarithm of the number of items stored in the tree.
- On average, BST with n nodes has O(log₂n) height.
- In the worst case, BST can have O(n) height.

Example



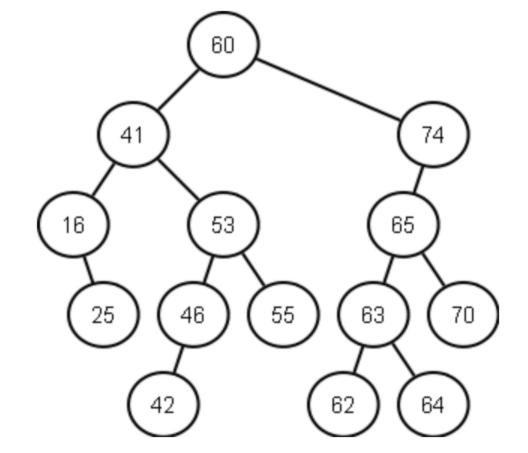
Structure code of a tree node

```
struct node
{ int data; //Data element
 struct node * left; //Pointer to left node
 struct node * right; //Pointer to right node
};
struct node * root = NULL;
```

Contd...

Traversals

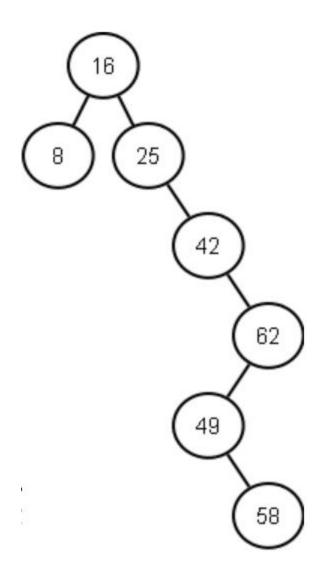
- Preorder traversal: (parent, left, right)
- Inorder traversal: (left, parent, right)
- Postorder traversal: (left, right, parent)

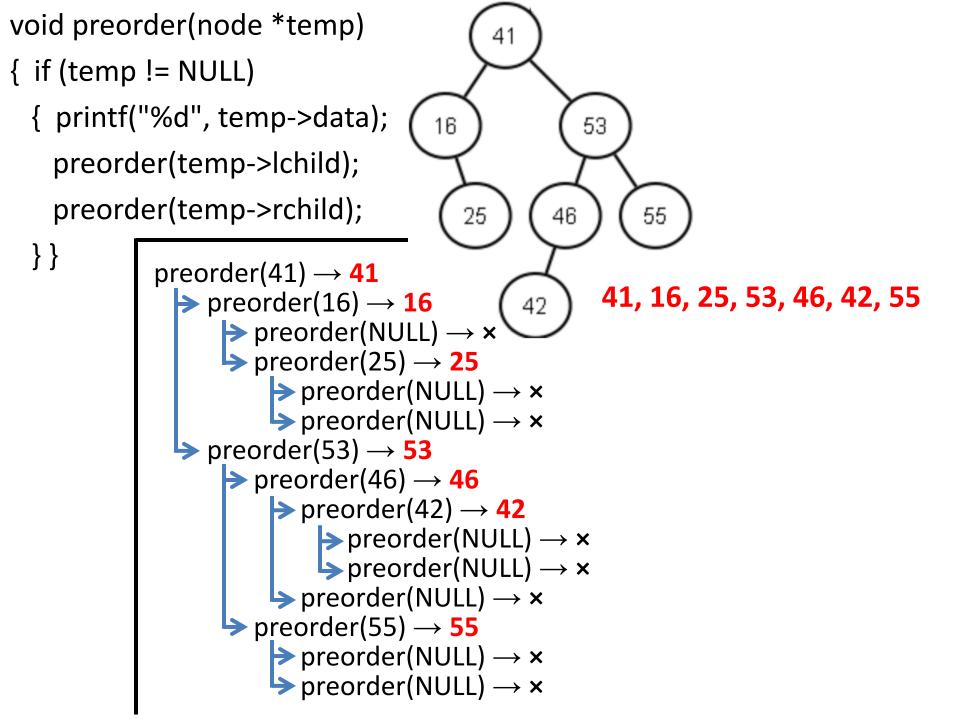


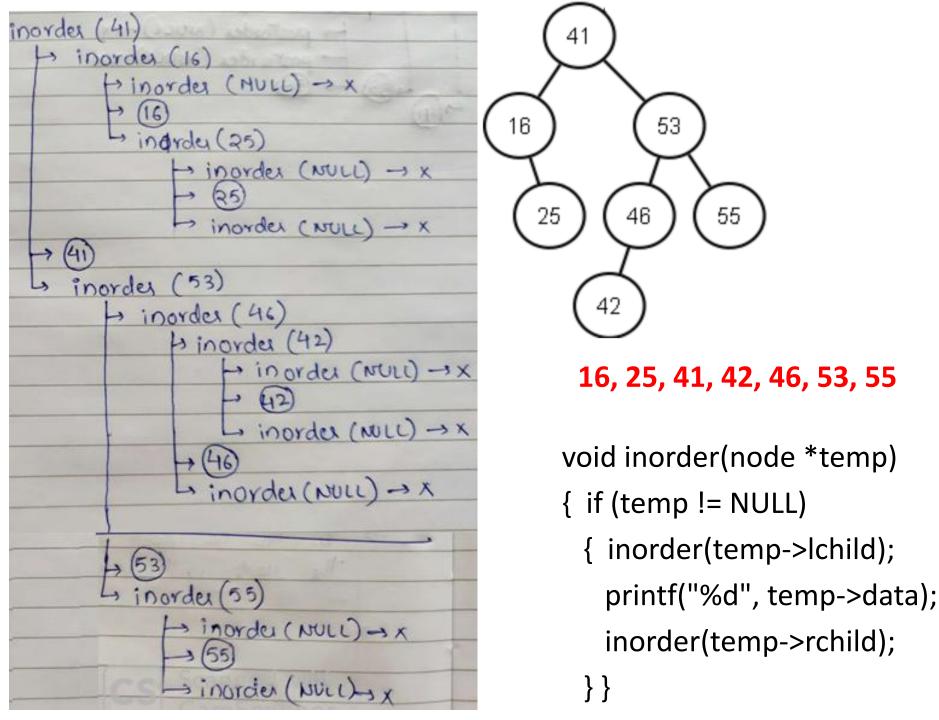
- Preorder traversal
 60, 41, 16, 25, 53, 46, 42, 55, 74, 65, 63, 62, 64, 70
- Inorder traversal
 16, 25, 41, 42, 46, 53, 55, 60, 62, 63, 64, 65, 70, 74
- Postorder traversal
 25, 16, 42, 46, 55, 53, 41, 62, 64, 63, 70, 65, 74, 60

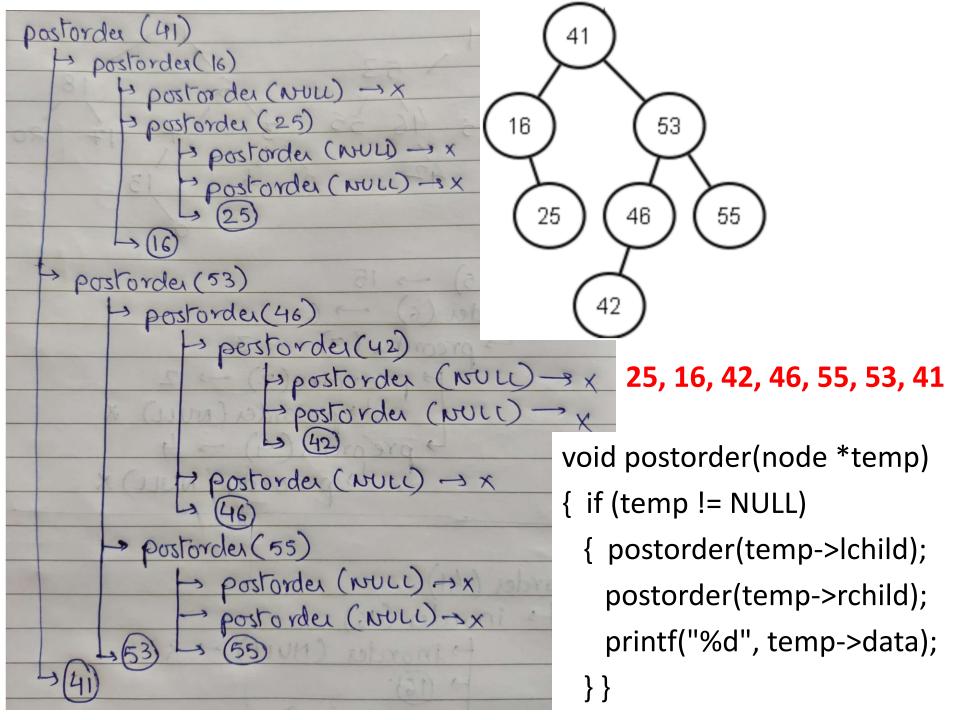
Traversals

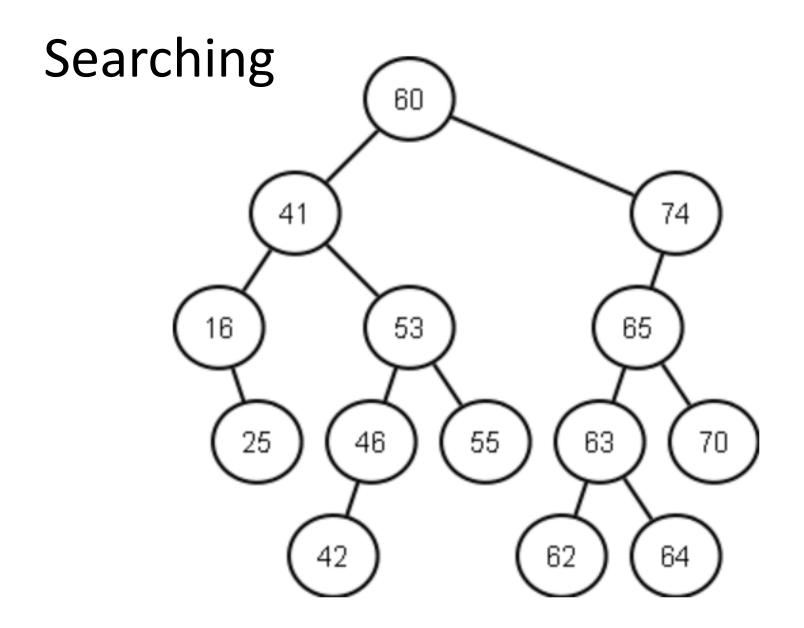
- Preorder traversal
 16, 8, 25, 42, 62, 49, 58
- Inorder traversal
 8, 16, 25, 42, 49, 58, 62
- Postorder traversal
 8, 58, 49, 62, 42, 25, 16











Contd...

```
TREE-SEARCH(x,k)

1 if x == NIL or k == x.key

2 return x

3 if k < x.key

4 return TREE-SEARCH(x.left,k)

5 else return TREE-SEARCH(x.right,k)
```

Running time is O(h), where h is the height of the tree.

```
ITERATIVE-TREE-SEARCH (x, k)

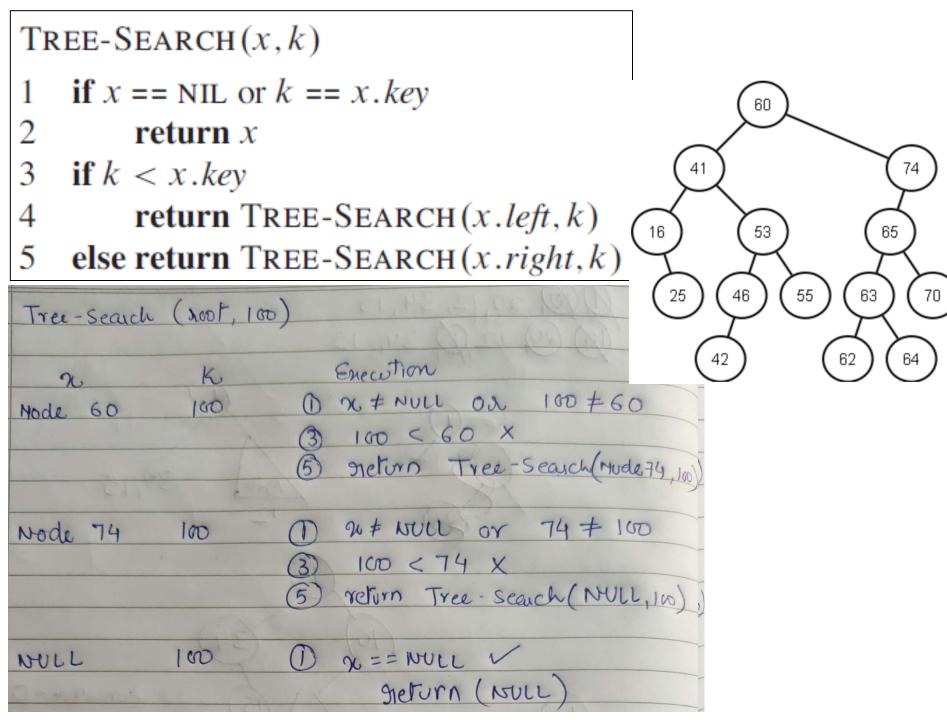
1 while x \neq \text{NIL} and k \neq x.key

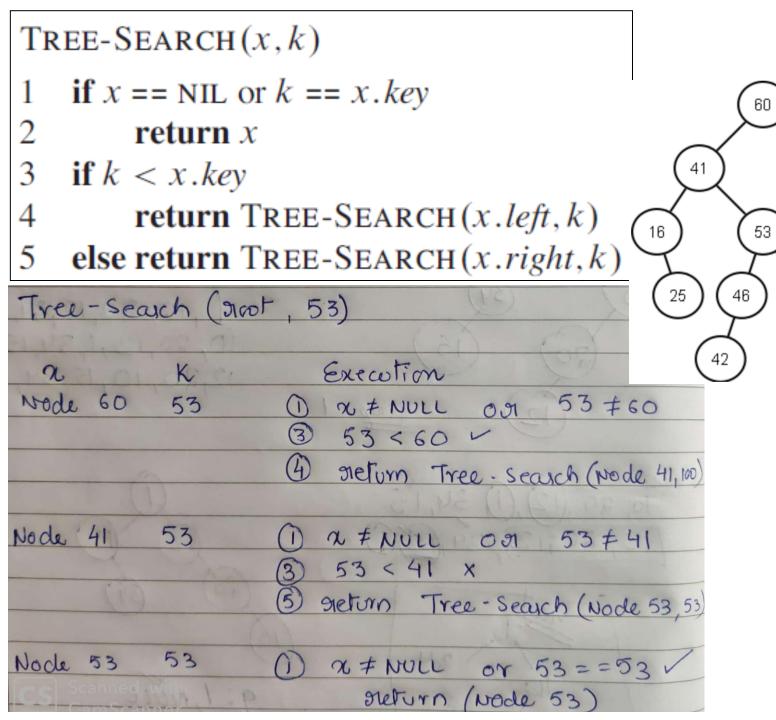
2 if k < x.key

3 x = x.left

4 else x = x.right

5 return x
```





```
ITERATIVE-TREE-SEARCH(x,k)

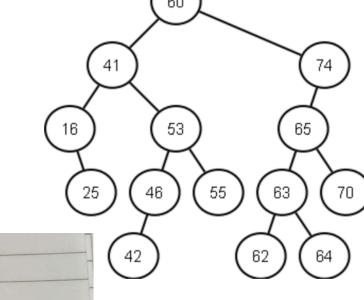
1 while x \neq \text{NIL} and k \neq x.key

2 if k < x.key

3 x = x.left

4 else x = x.right

5 return x
```



I terative -	Tree - Sec	arch (9100t, 100)
20	K	Execution
Node 60	100	91 \$ NULL V 100 \$ 60 V
		100 < 60 X
	100	n = n. right = Node 74
Node 74	100	91 \$ NULL V 100 \$ 74 V
		100 < 74 ×
		n = n. right = vode (vou)
NULL	100	20 \$ NULL X STEPTOM (NULL)

```
ITERATIVE-TREE-SEARCH (x, k)

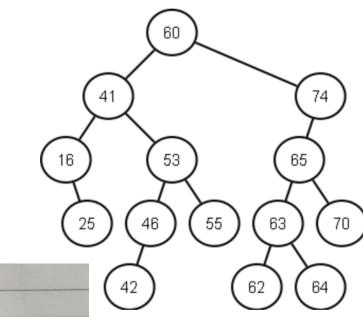
1 while x \neq \text{NIL} and k \neq x.key

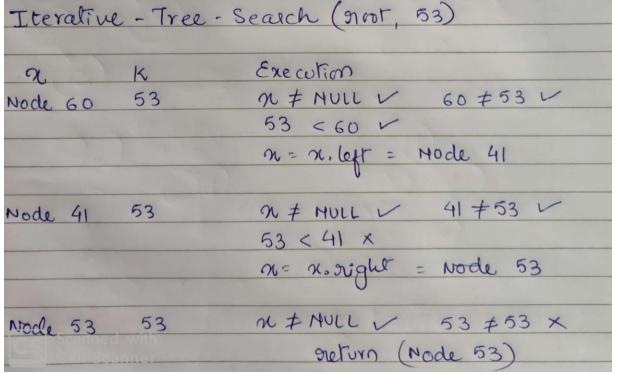
2 if k < x.key

3 x = x.left

4 else x = x.right

5 return x
```

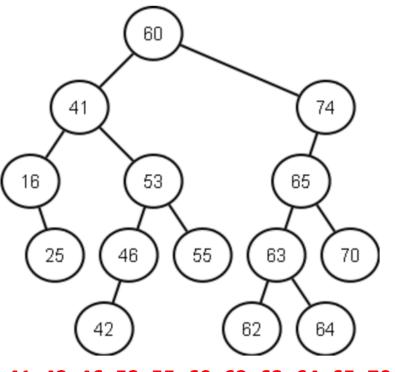




Minimum and Maximum

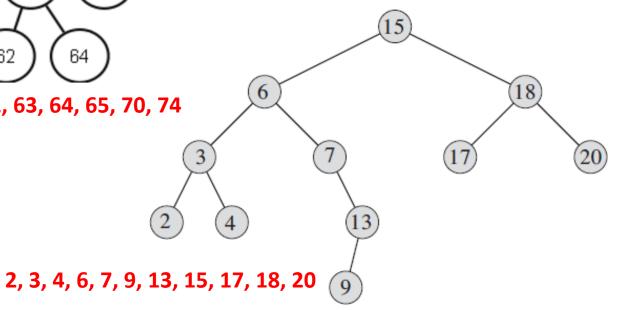
TREE-MAXIMUM(x) TREE-MINIMUM (x)while $x.right \neq NIL$ while $x.left \neq NIL$ x = x.leftx = x.rightreturn x return x Maximum (rightmost) Minimum (leftmost)

Successor and Predecessor



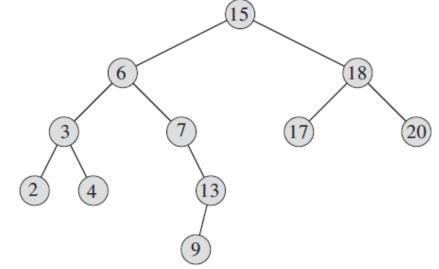
- Successor is next node in inorder traversal.
- Predecessor is previous node in inorder traversal.





TREE-SUCCESSOR (x)

- 1 **if** $x.right \neq NIL$
- 2 **return** TREE-MINIMUM (x.right)
- y = x.p
- 4 while $y \neq NIL$ and x == y.right
- $5 \qquad x = y$
- 6 y = y.p
- 7 **return** y

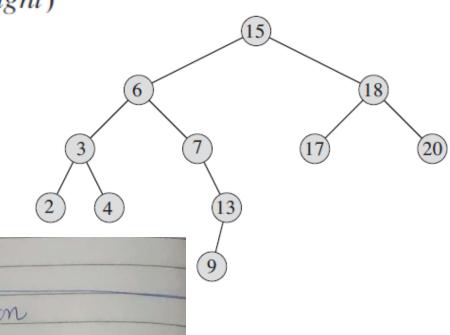


Tree.	-successor (6)		
76	n. right	7	Execution
Node 6	Node 7		1) Node 7 # NULL La Tree Minimum (Node7)
			L, (7)

TREE-SUCCESSOR (x)if $x.right \neq NIL$ **return** TREE-MINIMUM (x.right)(15)y = x.pwhile $y \neq NIL$ and x == y.right(18)x = yy = y.preturn y Tree - Successor (13) Execution n. right 90 1) NULL # NULL X Nocle 7 Node 13 NULL (4) Node 7 + NULL and Node 13 = = Node 13 Nocle 6 + NUIL and Node 6 Node 7 Node 7 == Noole 7 9 vode 15 + vull and Node 15 Node 6 Node 6 == Node 18X 4 (15)

TREE-SUCCESSOR (x)

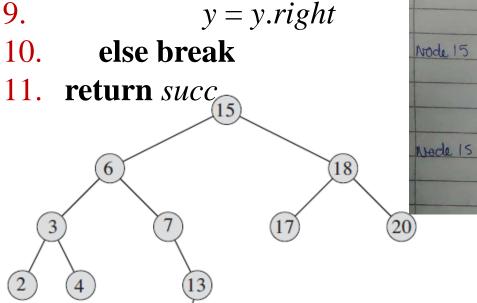
- 1 **if** $x.right \neq NIL$
- 2 **return** TREE-MINIMUM (x.right)
- y = x.p
- 4 while $y \neq NIL$ and x == y.right
- $5 \qquad x = y$
- 6 y = y.p
- 7 **return** y

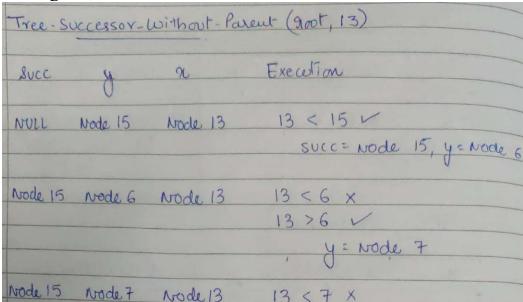


Tree - S	, uccessor ((20)	
20	n. right	T y	Execution
Node 20	NULL	Nocle 18	(4) Node 18 + NULL and
			Node 20 == Node 20
Node 18		Node 15	Node 18 == Nocle 18
CS to classifice	inher	NULL	1) NUCL # NOELX

Tree-Successor-Without-Parent(y,x)

- 1. **if** $x.right \neq NIL$
- 2. **return** Tree-Minimum(x.right)
- 3. succ = NIL
- 4. **while** TRUE
- $\mathbf{5.} \qquad \mathbf{if} \ x.key < y.key$
- 6. succ = y
- 7. y = y.left
- 8. **else if** x.key > y.key9. y = y.right





wede 13

Node 13

y= Node 13

break; neturn (15)

13 < 13 ×

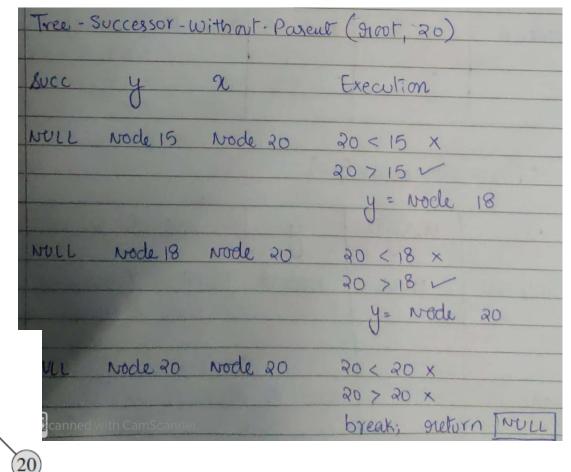
13 > 13 x

Tree-Successor-Without-Parent(y,x)

- 1. **if** $x.right \neq NIL$
- 2. **return** Tree-Minimum(x.right)

(18)

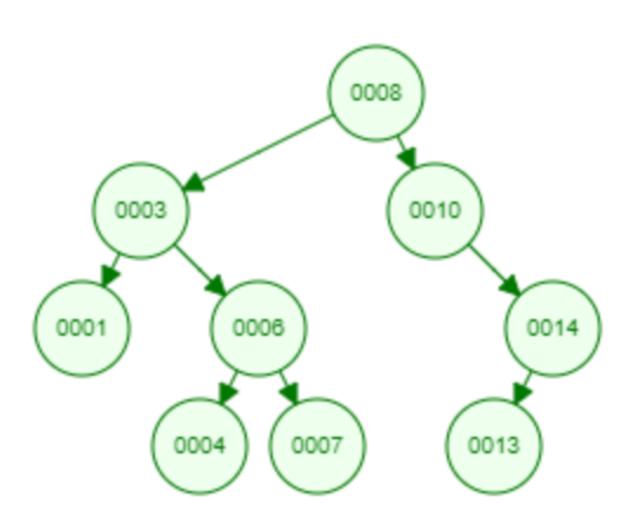
- 3. succ = NIL
- 4. **while** TRUE
- 5. **if** x.key < y.key
- 6. succ = y
- 7. y = y.left
- 8. else if x.key > y.key
 - y = y.right
- 10. else break
- 11. return succ



Insertion

- A new key is always inserted at leaf.
- Start searching a key from root till a leaf node is found.
 - Add the new node as a child of the leaf node.
- Duplicate key values are not allowed.
- If key to be inserted is already present, then return that node.

Create BST using 8, 3, 1, 10, 6, 14, 4, 7, 13



TREE-INSERT-WITHOUT-PARENT(T.root,z)

- 1. x = T.root
- 2. if x == NIL
- 3. return z
- 4. **if** z.key < x.key
- 5. x.left = Tree-Insert-Without-Parent(x.left,z)
- 6. **else if** z.key > x.key
- 7. x.right = Tree-Insert-Without-Parent(x.right,z)
- 8. return x

TREE-INSERT(T, z)y = NILx = T.root3 **while** $x \neq NIL$ 4 y = x5 if z.key < x.keyx = x.leftelse x = x.right $z \cdot p = y$ if y == NIL10 T.root = z // tree T was empty elseif z.key < y.keyy.left = zelse y.right = z

Contd...

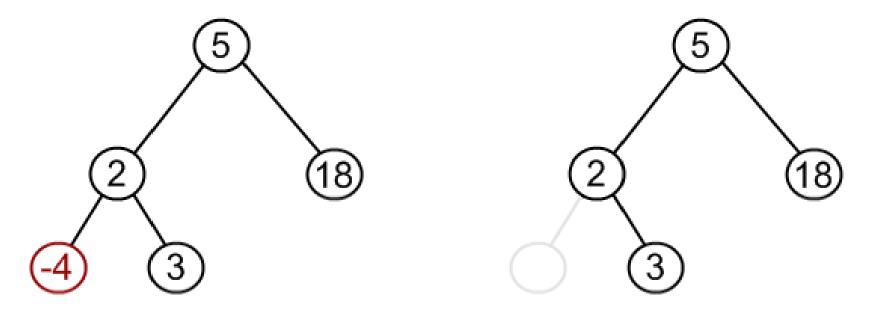
Deletion

- 1. Search for a node to remove;
- 2. If the node is found, execute the remove algorithm.

- Three cases,
 - Node to be removed has no children.
 - Node to be removed has one child.
 - iii. Node to be removed has two children.

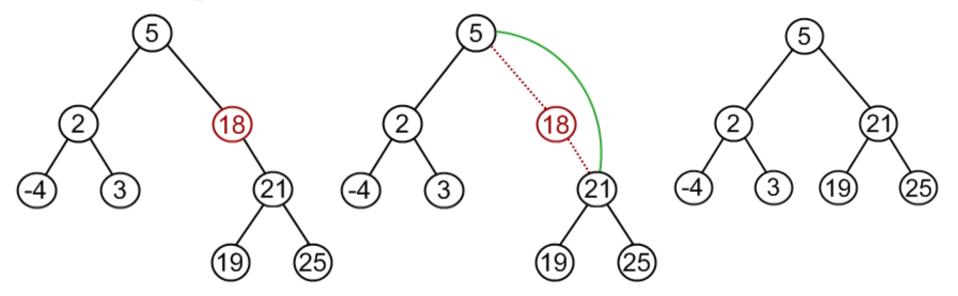
Node to be removed has no children

- Set corresponding link of the parent to NULL and disposes the node.
- Example. Remove -4 from a BST.



Node to be removed has one child

- Link single child (with it's subtree) directly to the parent of the removed node.
- Example. Remove 18 from a BST.

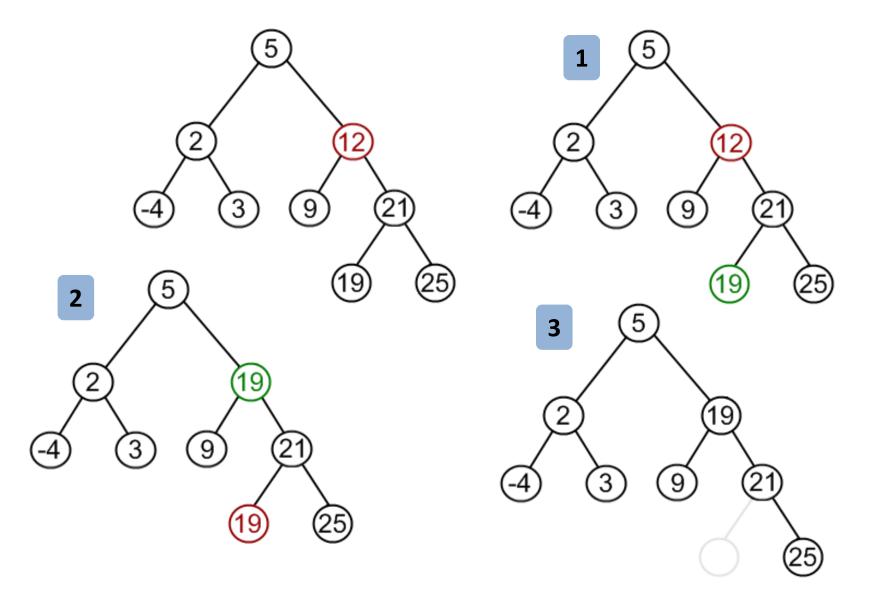


Node to be removed has two children

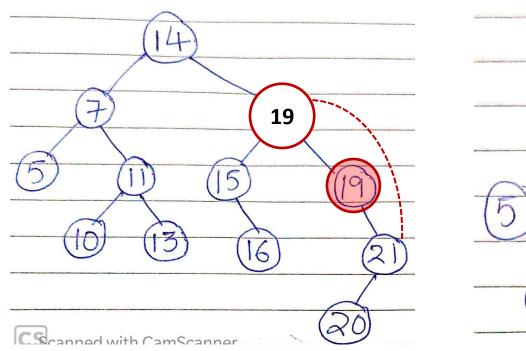
- Find inorder successor, i.e. find the minimum value in right child of the node.
- Copy contents of the inorder successor (or found minimum) to the node being removed with.
- Delete the inorder successor (or found minimum) from the right subtree.
- Note:
 - The node with minimum value has no left child and, therefore, it's removal may result in first or second cases only.
 - Inorder predecessor can also be used.

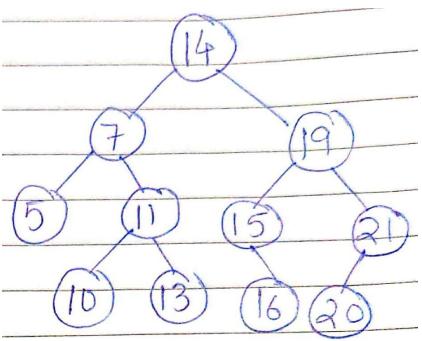
Contd...

Remove 12 from a BST.



Delete 17





```
TRANSPLANT (T, u, v)

1 if u.p == NIL

2 T.root = v

3 elseif u == u.p.left

4 u.p.left = v

5 else u.p.right = v

6 if v \neq NIL

7 v.p = u.p
```

```
Tree-Delete (T, z)
 1 if z.left == NIL
        TRANSPLANT(T, z, z.right)
 3 elseif z.right == NIL
        TRANSPLANT(T, z, z.left)
    else y = \text{Tree-Minimum}(z.right)
 6
        if y.p \neq z
             Transplant(T, y, y.right)
 8
             y.right = z.right
             y.right.p = y
        TRANSPLANT(T, z, y)
10
        y.left = z.left
11
        y.left.p = y
12
```

```
TREE-DELETE-WITHOUT-PARENT(T.root,k)
    x = T.root
2. 	 if x == NIL
3.
         return x
4.
    if k < x.key
5.
        x.left = Tree-Delete-Without-Parent(x.left,k)
    else if k > x.key
6.
7.
         x.right = \text{Tree-Delete-Without-Parent}(x.right,k)
8.
     else
9.
         if x.left == NIL
10.
                temp = x.right
11.
                 delete x
12.
                return temp
        else if x.right == NIL
13.
                 temp = x.left
14.
15.
                 delete x
16.
                 return temp
17.
        temp = TREE-MINIMUM(x.right)
18.
        x.key = temp.key
         x.right = Tree-Delete-Without-Parent(x.right,temp.key)
19.
20.
     return x
```

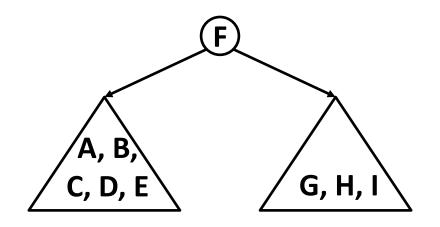
Question...

15, 18, 6, 7, 17, 3, 4, 13, 9, 20, 2

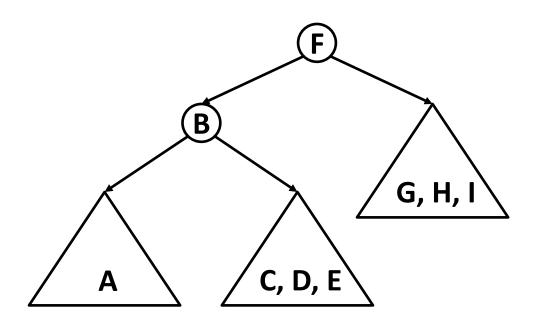
- Generate BST for the given sequence.
- Visit the generated BST using inorder, preorder, and postorder traversals.
- Delete 4, 7, and 15 in sequence showing BST after every deletion.

Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]

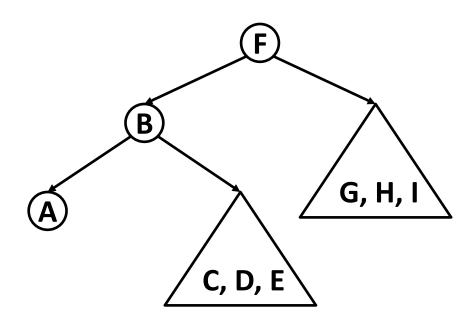
Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]



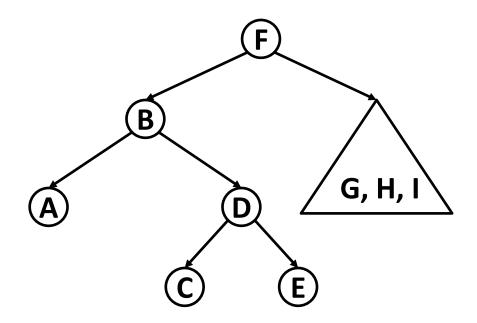
Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]



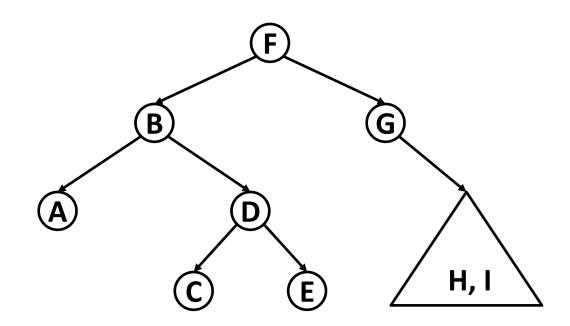
Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]



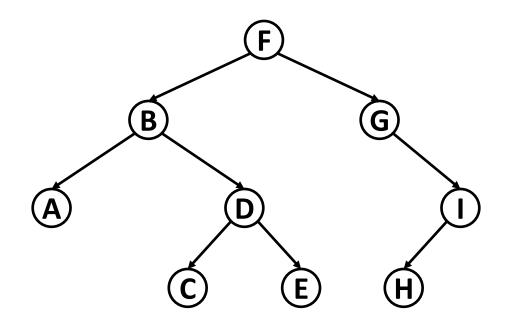
Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]



Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]

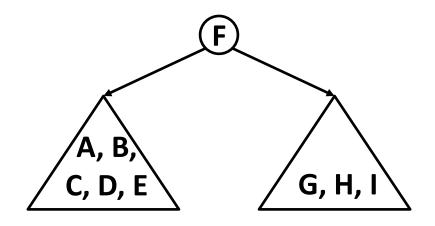


Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]

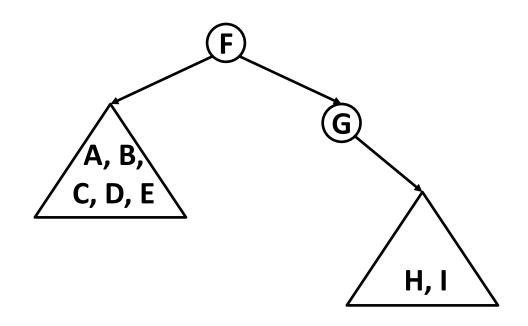


Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]

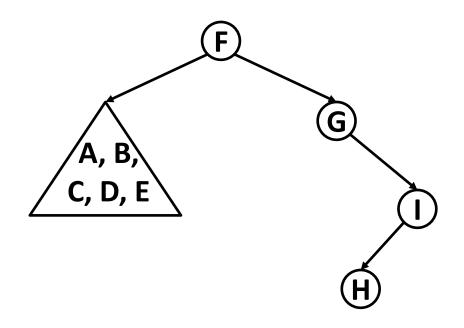
Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]



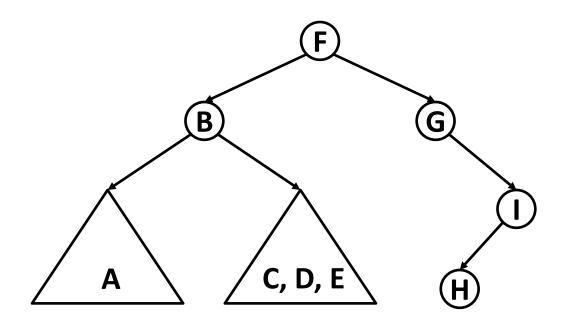
Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]



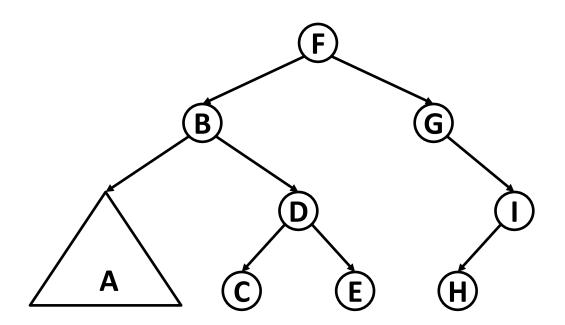
Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]



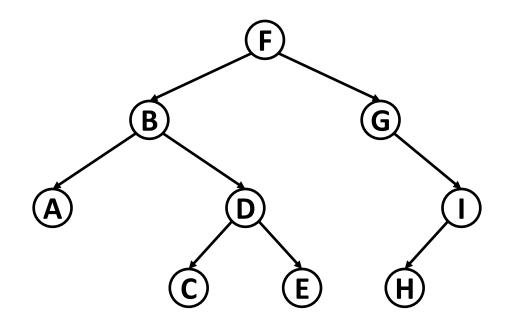
Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]



Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]

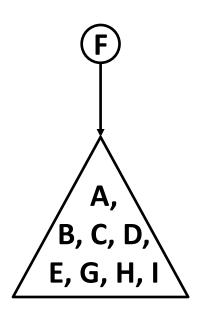


Inorder: A, B, C, D, E, F, G, H, I [LEFT PARENT RIGHT]

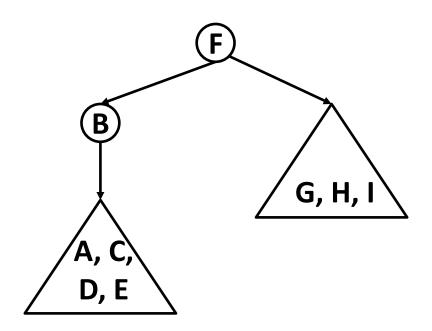


Preorder: F, B, A, D, C, E, G, I, H [PARENT LEFT RIGHT]

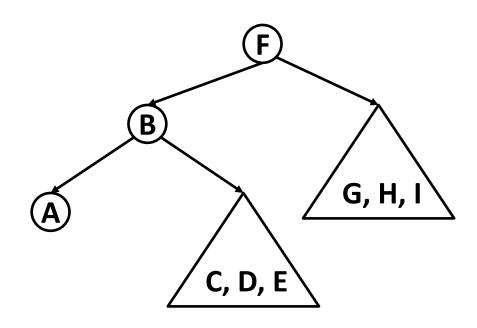
Preorder: F, B, A, D, C, E, G, I, H [PARENT LEFT RIGHT]



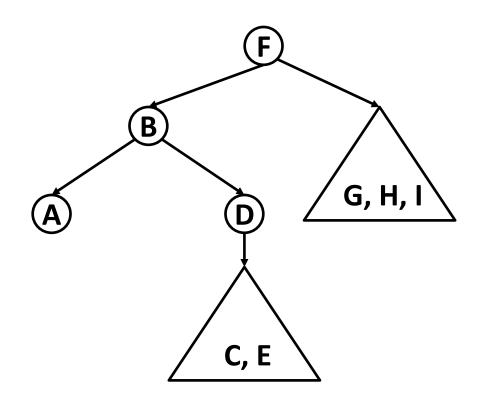
Preorder: F, B, A, D, C, E, G, I, H [PARENT LEFT RIGHT]



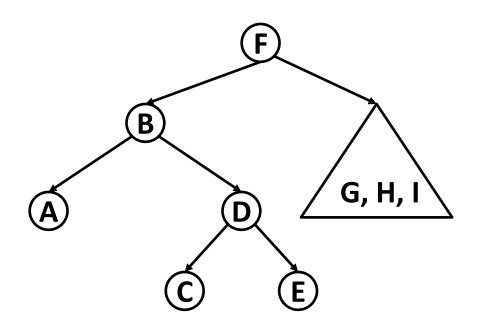
Preorder: F, B, A, D, C, E, G, I, H [PARENT LEFT RIGHT]



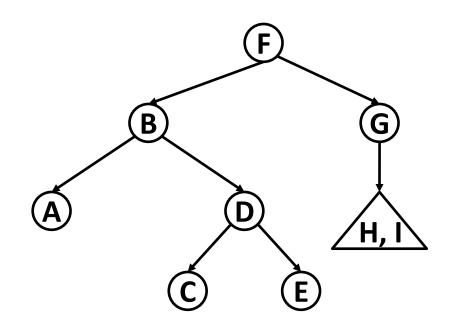
Preorder: F, B, A, D, C, E, G, I, H [PARENT LEFT RIGHT]



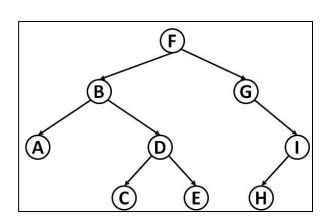
Preorder: F, B, A, D, C, E, G, I, H [PARENT LEFT RIGHT]

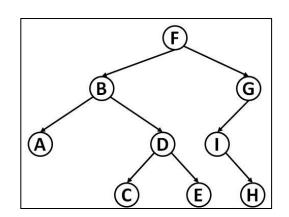


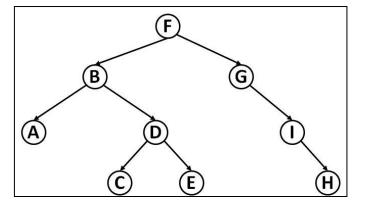
Preorder: F, B, A, D, C, E, G, I, H [PARENT LEFT RIGHT]

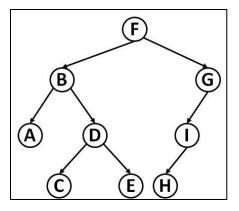


Preorder: F, B, A, D, C, E, G, I, H [PARENT LEFT RIGHT]



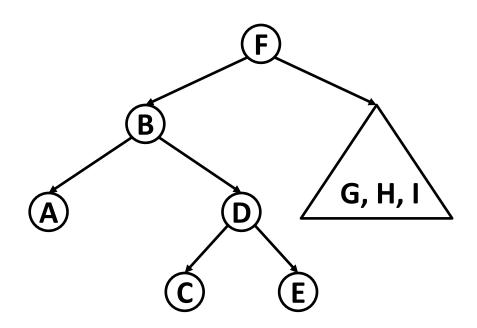






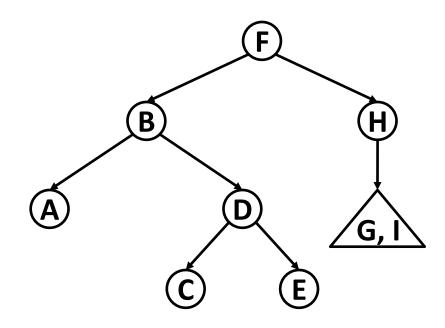
Preorder and Postorder (Full Tree)

Preorder: F, B, A, D, C, E, H, G, I [LEFT PARENT RIGHT]



Preorder and Postorder (Full Tree)

Preorder: F, B, A, D, C, E, H, G, I [LEFT PARENT RIGHT]



Preorder and Postorder (Full Tree)

Preorder: F, B, A, D, C, E, H, G, I [LEFT PARENT RIGHT]

