

Roll Number: _____

Number of Pages: 02

Thapar Institute of Engineering and Technology, Patiala

Department of Computer Science and Engineering

END SEMESTER EXAMINATION

B. E. (Second Year): Sem-I (2022/23) Course Code: UCS301 Course Name: Data Structures

December 12, 2022 Monday, 16:30 Hrs – 19:30 Hrs Time: 3 Hours, M. Marks: 40

Name of Faculty: SP, RMT, RKR, VBP, SMT, Dr. Manisha, Dr. Diwakar

Note: Attempt subparts of a question in sequence at one place. Assume missing data, if any, suitably.

- Q1. (a) The height of a tree is the length of the longest path from root to leaf. Find the height of the binary tree whose respective postorder and inorder traversals are 8, 9, 6, 7, 4, 5, 2, 3, 1 and 8, 6, 9, 4, 7, 2, 5, 1, 3. (4 marks)

- (b) Convert the BST shown in Fig. 1. into an AVL tree, showing trees at each intermediate step. (2 marks)

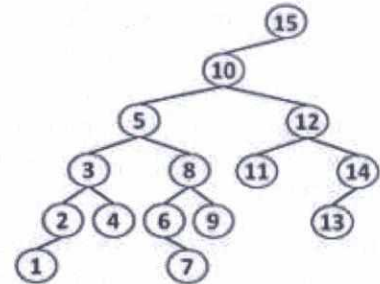


Fig. 1

- Q2. (a) For the graph shown in Fig. 2, determine the sequence in which edges get selected by Prim's minimum spanning tree algorithm. Start from vertex A. (3 marks)

- (b) Execute Dijkstra's single source shortest path algorithm on the graph given in Fig. 3. Consider vertex T as the starting vertex. (4 marks)

	A	B	C	D	E	F	G
A	0	7	0	5	0	0	0
B	7	0	8	9	7	0	0
C	0	8	0	0	5	0	0
D	5	9	0	0	15	6	0
E	0	7	5	15	0	8	9
F	0	0	0	6	8	0	11
G	0	0	0	0	9	11	0

Fig. 2

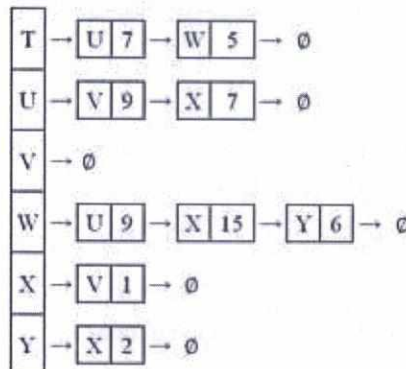


Fig. 3

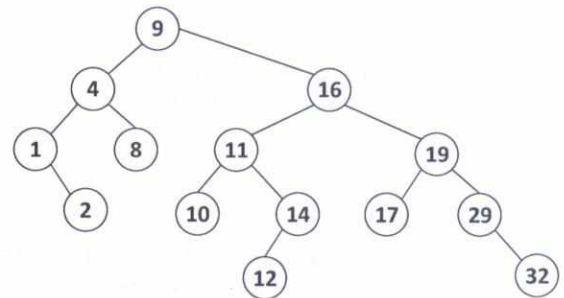


Fig. 4

- Q3. (a) Build an AVL tree by sequentially inserting the keys (45, 20, 32, 30, 55, 65, 25, 40, 29, 35, and 50). After each insertion that includes a rotation, redraw the tree. (5 marks)

- (b) Delete the node 8 from the AVL tree shown in Fig. 4 and draw the resulting AVL tree. (2 marks)

BUILD-MAX-HEAP(A)

```

1  A.heap-size = A.length
2  for i = [A.length/2] downto 1
3      MAX-HEAPIFY(A, i)

```

MAX-HEAPIFY(A, i)

```

1  l = LEFT(i)
2  r = RIGHT(i)
3  if l ≤ A.heap-size and A[l] > A[i]
4      largest = l
5  else largest = i
6  if r ≤ A.heap-size and A[r] > A[largest]
7      largest = r
8  if largest ≠ i
9      exchange A[i] with A[largest]
10 MAX-HEAPIFY(A, largest)

```

Fig. 5

- Q4. Execute BUILD-MAX-HEAP (Fig. 5) on the array $A[] = \{5, 3, 17, 10, 84, 19, 6, 22, 9\}$. Show all intermediate steps involved in the construction of MAX-HEAP. Also prove that the time complexity of BUILD-MAX-HEAP is $O(n)$, where n is the total number of elements in the heap. (3 marks) + (3 marks)

Q5. (a) Consider an unsorted array of **13** integers. You have asked to sort it using merge-sort algorithm. Find out the total number of **worst-case** and **best-case** comparisons needed to run the merge sort algorithm? Justify your answers by giving proper example. (Hint: The total number of comparisons is the exact number of comparisons, not the approximate comparisons.) (4 marks)

(b) You have given an unsorted array having elements **25, 47, 15, 8, 9, 4, 40, 30, 10, 19**, and asked to sort the array using quick sort algorithm. You select the last element **19** as the pivot element. Show the content of the left sub-array and right sub-array after fixing **19** at its correct index assuming that the array starts with the index **zero**. No need to sort the entire array. (3 marks)

Q6. (a) Consider a hash table shown in Fig. 6. Keys are inserted using double hashing collision resolution technique with $h(k, i) = (h_1(k) + ih_2(k)) \bmod m$, $m = 11$, $h_1(k) = k$, and $h_2(k) = 1 + k \bmod (m - 1)$. Determine the correct indices giving the complete probe sequence for keys **22** and **31**. (2 marks)

0	1	2	3	4	5	6	7	8	9	10
88		28		59		17			4	15

Fig. 6

(b) Execute **function_one (node *root)** in Fig. 7 on the BST shown in Fig. 1 and explain the purpose of designing it. (Hint: Try to draw trees obtained in intermediate steps of the algorithm.) (5 marks)

```

struct node { int key;   node *left, *right; };

int height (node *root); // Returns zero for a single node tree.
node* rightRotate (node *root); // Performs right rotation.
node* leftRotate (node *root); // Performs left rotation.

node* function_two (node *root)
{ // Declare two node pointers p and r. Allocate memory dynamically to pointer r.
  r -> left = NULL;
  r -> right = root;
  p = r;
  while (p -> right) {
    if (p -> right -> left == NULL)    p = p -> right;
    else p -> right = rightRotate(p -> right);
  }
  return r;
}

node* function_three (node *root, int count )
{ // Declare a node pointer p and initialize it to root.
  while (count) {
    p -> right = leftRotate(p -> right);
    p = p -> right;
    --count;
  }
  return root;
}

node* function_one (node *root)
{ root = function_two(root);
  // Declare integers n, t, l, and lc.
  n = height(T); // T = root
  t = n + 1; l = 0;
  while (t > 1) { ++l; t /= 2; }
  lc = n + 1 - 2^l;
  if (lc == 0)    lc = 2^(l - 1);
  root = function_three(root,lc);
  n -= lc;
  while (n > 1) {
    n /= 2;
    root = function_three(root,n);
  }
  return T; // T = root
}

```

Fig. 7

-----ALL THE BEST-----