Recurrences

Instructor: Dr. Tarunpreet Bhatia
Assistant Professor, CSED
Thapar Institute of Engineering and Technology

Recurrence Relation

- Many algorithms (such as divide and conquer) are recursive in nature.
- When an algorithm contains a recursive call to itself then its running time can be described by mathematical relation called recurrence relation or recurrence equation.
- Recurrence relations are used to determine the running time of a recursive program.

Recurrence Relation

- We get running time as a function of *n* (input size) and we get the running time on inputs of smaller sizes.
- A recurrence is a recursive description of a function, or a description of a function in terms of itself.
- A recurrence relation recursively defines a sequence where the next term is a function of the previous terms.

$$s(n) = \begin{cases} 0 & n = 0 \\ c + s(n-1) & n > 0 \end{cases} \qquad s(n) = \begin{cases} 0 & n = 0 \\ n + s(n-1) & n > 0 \end{cases}$$

$$s(n) = \begin{cases} 0 & n = 0\\ n + s(n-1) & n > 0 \end{cases}$$

$$T(n) = \begin{cases} c & n = 1 \\ 2T\left(\frac{n}{2}\right) + c & n > 1 \end{cases}$$

$$T(n) = \begin{cases} c & n = 1 \\ 2T\left(\frac{n}{2}\right) + c & n > 1 \end{cases}$$

$$T(n) = \begin{cases} c & n = 1 \\ aT\left(\frac{n}{b}\right) + cn & n > 1 \end{cases}$$

```
void display (int n)
{
    if (n > 0)
    {
       printf ("%d", n);
       display (n -1);
    }
}
```

Recurrence Modified Example 1

```
void display (int n)
{
    if (n > 1)
    {
       printf ("%d", n);
       display (n -1);
    }
}
```

```
void test (int n)
{
    if (n > 0)
    {
        printf("%d",n);
        test (n-1);
        test (n-1);
    }
}
```

```
void test (int n)
{
    if (n > 0)
    {
        printf("%d", n);
        test (n/2);
    }
}
```

Example: Binary search using recursion

```
int binarySearch(int arr[], int l, int r, int x)
      if (l<=r) {
      int mid = 1 + (r - 1) / 2;
      if (arr[mid] == x)
         return mid;
      if (arr[mid] > x)
         return binarySearch(arr, l, mid - 1, x);
      return binarySearch(arr, mid + 1, r, x);
      return -1;
```

```
void test (int n)
{
    if (n > 0)
    {
        printf("%d", n);
        test (n/2);
        test (n/2);
    }
}
```

```
void test (int n)
{
    if (n > 1)
    {
       test (2n/3);
       test (n/3);
    }
}
```

Recursion Example 8

```
\label{eq:continuous_section} $$\inf n < 2$$ return n; $$ return doSomething(n-1) + doSomething(n-2); $$$
```

Solving Recurrences

- Iteration method or Back substitution
- Recursion Tree method
- Master method

Iteration Method

- In iteration method,
 - Expand the recurrence
 - Work some algebra to express as a summation
 - Evaluate the summation
- We will show several examples

Recurrence Example 1 $T(n) = \begin{cases} 1 & n = 0 \\ 1 + T(n-1) & n > 0 \end{cases}$

Recurrence Example 2
$$T(n) = \begin{cases} 1 & n = 0 \\ n + T(n-1) & n > 0 \end{cases}$$

$$T(n) = \begin{cases} 1 & n = 0 \\ T(n-1) + \log n & n > 0 \end{cases}$$

$$T(n) = \begin{cases} 1 & n = 0 \\ 2T(n-1) + 1 & n > 0 \end{cases}$$

$$T(n) = \begin{cases} 1 & n = 1 \\ T(n/2) + 1 & n > 1 \end{cases}$$

$$T(n) = \begin{cases} 1 & n = 1 \\ 2T(n/2) + 1 & n > 1 \end{cases}$$

Examples

$$T(n) = T(n-1) + 1 - - - - O(n)$$

$$T(n) = T(n-1) + n - - - - O(n^{2})$$

$$T(n) = T(n-1) + n^{2} - - - - O(n^{3})$$

$$T(n) = T(n-1) + \log n - - - O(n \log n)$$

$$T(n) = T(n-2) + 1 \rightarrow \frac{n}{2} \rightarrow O(n)$$

$$T(n) = T(n-50) + n - - - - O(n^{2})$$

Try this!

$$T(n) = \begin{cases} 1 & n=1 \\ T(n/2) + n & n>1 \end{cases}$$

$$T(n) = \begin{cases} 1 & n=1\\ 2T((n/2)+16)+n & n>1 \end{cases}$$

$$T(n) = \begin{cases} 1 & n=1 \\ T(n/3) + n^4/3 & n>1 \end{cases}$$

Recurrence Relation for Root Function

Recurrence Relation: $T(n) = T(\sqrt{n}) + 1$

$$T(n) = \begin{cases} 1 & n = 2 \\ T(\sqrt{n}) + 1 & n > 2 \end{cases}$$

$$T(n) = T(\sqrt{n}) + 1$$

$$T(n) = T(n^{\frac{1}{2}}) + 1$$

$$T(n) = T(n^{\frac{1}{2^2}}) + 2$$

$$T(n) = T(n^{\frac{1}{2^3}}) + 3$$

 $T(n) = T(n^{\frac{1}{2^k}}) + k$

Assume
$$n = 2^m$$

$$T(2^m) = T(2^{\frac{m}{2^k}}) + k$$
Assume $T(2^{\frac{m}{2^k}}) = T(2^1)$
Therefore, $\frac{m}{2^k} = 1$

$$m = 2^k \text{ and } k = \log_2 m$$
Since $n = 2^m$, $m = \log_2 n$

$$k = \log \log_2 n$$

$$\theta(\log \log_2 n)$$

$$T(n) = 2T(\sqrt{n}) + \log n$$

Try this!!

$$T(n) = 2T(\sqrt{n}) + 1$$

$$T(n) = T(n-1) + (1/n)$$

$$T(n) = \sqrt{n} T(\sqrt{n}) + n$$

$$T(n) = 5T\left(\frac{n}{5}\right) + \frac{n}{\log n}$$
 where $T(1) = 1$