

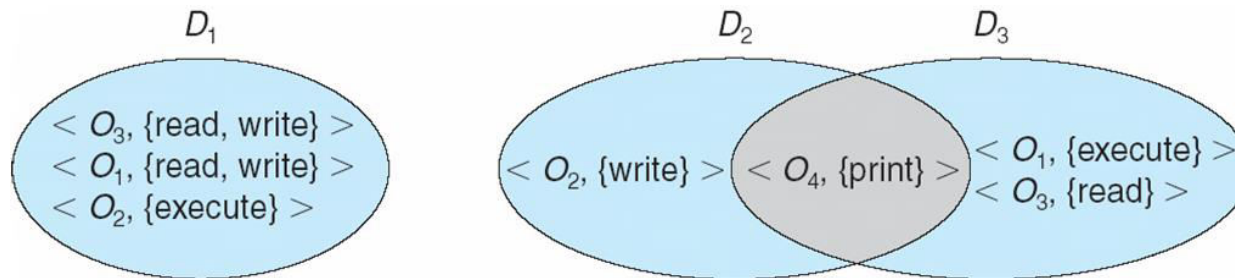
Chapter 14: Protection





Domain Structure

- Access-right = $\langle \text{object-name}, \text{rights-set} \rangle$
where *rights-set* is a subset of all valid operations that can be performed on the object
- Domain = set of access-rights





Access Matrix

- View protection as a matrix (**access matrix**)
- Rows represent domains
- Columns represent objects
- **Access**(i, j) is the set of operations that a process executing in Domain $_i$ can invoke on Object $_j$

domain \ object	F_1	F_2	F_3	printer
D_1	read		read	
D_2				print
D_3		read	execute	
D_4	read write		read write	





Use of Access Matrix

- If a process in Domain D_i tries to do “op” on object O_j , then “op” must be in the access matrix
- User who creates object can define access column for that object
- Can be expanded to dynamic protection
 - Operations to add, delete access rights
 - Special access rights:
 - ▶ *owner of O_i*
 - ▶ *copy op from O_i to O_j (denoted by “*”)*
 - ▶ *control – D_i can modify D_j access rights*
 - ▶ *transfer – switch from domain D_i to D_j*
 - *Copy and Owner* applicable to an object
 - *Control* applicable to domain object





Use of Access Matrix (Cont.)

- **Access matrix** design separates mechanism from policy
 - Mechanism
 - ▶ Operating system provides access-matrix + rules
 - ▶ If ensures that the matrix is only manipulated by authorized agents and that rules are strictly enforced
 - Policy
 - ▶ User dictates policy
 - ▶ Who can access what object and in what mode
- But doesn't solve the general confinement problem





Access Matrix of Figure A with Domains as Objects

domain \ object	F_1	F_2	F_3	laser printer	D_1	D_2	D_3	D_4
D_1	read		read			switch		
D_2				print			switch	switch
D_3		read	execute					
D_4	read write		read write		switch			





Access Matrix with Copy Rights

object domain	F_1	F_2	F_3
D_1	execute		write*
D_2	execute	read*	execute
D_3	execute		

(a)

object domain	F_1	F_2	F_3
D_1	execute		write*
D_2	execute	read*	execute
D_3	execute	read	

(b)





Access Matrix With Owner Rights

object domain	F_1	F_2	F_3
D_1	owner execute		write
D_2		read* owner	read* owner write
D_3	execute		

(a)

object domain	F_1	F_2	F_3
D_1	owner execute		write
D_2		owner read* write*	read* owner write
D_3		write	write

(b)





Modified Access Matrix of Figure B

domain \ object	F_1	F_2	F_3	laser printer	D_1	D_2	D_3	D_4
D_1	read		read			switch		
D_2				print			switch	switch control
D_3		read	execute					
D_4	write		write		switch			





Implementation of Access Matrix

- Generally, a sparse matrix
- Option 1 – Global table
 - Store ordered triples **<domain, object, rights-set>** in table
 - A requested operation M on object O_j within domain D_i -> search table for $\langle D_i, O_j, R_k \rangle$
 - ▶ with $M \in R_k$
 - But table could be large -> won't fit in main memory
 - Difficult to group objects (consider an object that all domains can read)





Implementation of Access Matrix (Cont.)

- Option 2 – Access lists for objects
 - Each column implemented as an access list for one object
 - Resulting per-object list consists of ordered pairs **<domain, rights-set>** defining all domains with non-empty set of access rights for the object
 - Easily extended to contain default set -> If $M \in \text{default set}$, also allow access





Implementation of Access Matrix (Cont.)

- Each column = Access-control list for one object
Defines who can perform what operation

Domain 1 = Read, Write
Domain 2 = Read
Domain 3 = Read

- Each Row = Capability List (like a key)
For each domain, what operations allowed on what objects
 - Object F1 – Read
 - Object F4 – Read, Write, Execute
 - Object F5 – Read, Write, Delete, Copy





Implementation of Access Matrix (Cont.)

- Option 3 – Capability list for domains
 - Instead of object-based, list is domain based
 - **Capability list** for domain is list of objects together with operations allows on them
 - Object represented by its name or address, called a **capability**
 - Execute operation M on object O_j , process requests operation and specifies capability as parameter
 - ▶ Possession of capability means access is allowed
 - Capability list associated with domain but never directly accessible by domain
 - ▶ Rather, protected object, maintained by OS and accessed indirectly
 - ▶ Like a “secure pointer”
 - ▶ Idea can be extended up to applications





Implementation of Access Matrix (Cont.)

- Option 4 – Lock-key
 - Compromise between access lists and capability lists
 - Each object has list of unique bit patterns, called **locks**
 - Each domain as list of unique bit patterns called **keys**
 - Process in a domain can only access object if domain has key that matches one of the locks

