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4 Matching & Flow 4.1 Kuhn Munkres	<pre>1.2 Increase Stack //stack resize(change esp to rsp if 64-bit system) asm( "mov %0,%%esp\n" ::"g"(mem+10000000) ); // craziest way static void run_stack_sz(void(*func)(),size_t stsize){ char *stack, *send;</pre>
5 Math       12         5.1 Prime Table       12         5.2 ⌊ ½ ⌋ Enumeration       12         5.3 ax+by=gcd       12         5.4 Pollard Rho       12         5.5 Pi Count (Linear Sieve)       12         5.6 Range Sieve       13         5.7 Miller Rabin       13         5.8 Inverse Element       13         5.9 Euler Phi Function       13         5.10Gauss Elimination       13         5.11Fast Fourier Transform       13         5.12High Speed Linear Recurrence       14         5.13Chinese Remainder       14         5.14Berlekamp Massey       15         5.15NTT       15         5.16Polynomial Operations       15         5.17FWT       16         5.18DiscreteLog       16         5.19Quadratic residue       16         5.20De-Bruijn       17         5.21Simplex Construction       17         5.22Simplex       17	<pre>send=(char *)((uintptr_t)send/16*16); asm volatile(    "mov %%rsp, (%0)\n"    "mov %0, %%rsp\n"    :    : "r" (send)); func(); asm volatile(    "mov (%0), %%rsp\n"    :    : "r" (send)); free(stack); }  1.3 Pragma optimization  // #pragma GCC optimize("Ofast, no-stack-protector") #pragma GCC optimize("Dfast, no-stack-protector") #pragma GCC optimize("Dfast, no-stack-protector")</pre>
6 Geometry 6.1 Point Class	<pre>#pragma GCC target("popcnt,abm,mmx,avx,tune=native")  1.4 IO Optimization  static inline int gc() {     static char buf[ 1 &lt;&lt; 20 ], *p = buf, *end = buf;     if ( p == end ) {         end = buf + fread( buf, 1, 1 &lt;&lt; 20, stdin );         if ( end == buf ) return EOF;         p = buf; }</pre>
7 Stringology 21 7.1 Hash	<pre>} template &lt; typename T &gt; static inline bool gn( T &amp;_ ) {     register int c = gc(); register T = 1; _ = 0;     while(('0'&gt;c  c&gt;'9') &amp;&amp; c!=EOF &amp;&amp; c!='-') c = gc();     if(c == '-') { = -1; c = gc(); }     if(c == EOF) return false;     while('0'&lt;=c&amp;&amp;c&lt;='9') _ = _ * 10 + c - '0', c = gc();     _ *=;</pre>

```
}
template < typename T, typename ...Args >
static inline bool gn( T &x, Args &...args )
{ return gn(x) && gn(args...); }
```

## 2 Data Structure

## 2.1 Bigint

```
class BigInt{
private
  using lld = int_fast64_t;
  #define PRINTF_ARG PRIdFAST64
  #define LOG_BASE_STR "9"
  static constexpr lld BASE = 10000000000;
  static constexpr int LOG_BASE = 9;
  vector<lld> dig;
 bool neg;
  inline int len() const { return (int) dig.size(); }
  inline int cmp_minus(const BigInt& a) const {
    if(len() == 0 && a.len() == 0) return 0;
    if(neg`^ a.neg)return (int)a.neg*2 - 1;
    if(len()!=a.len())
    return neg?a.len()-len():len()-a.len();
for(int i=len()-1;i>=0;i--) if(dig[i]!=a.dig[i])
      return neg?a.dig[i]-dig[i]:dig[i]-a.dig[i];
    return 0;
  inline void trim(){
    while(!dig.empty()&&!dig.back())dig.pop_back();
    if(dig.empty()) neg = false;
public:
 BigInt(): dig(vector<lld>()), neg(false){}
 BigInt(lld a): dig(vector<lld>()){
    neg = a<0; dig.push_back(abs(a));</pre>
    trim();
  BigInt(const string& a): dig(vector<lld>()){
    assert(!a.empty()); neg = (a[0]=='-');
    for(int i=((int)a.size())-1;i>=neg;i-=LOG_BASE){
      11d cur = 0;
      for(int j=min(LOG_BASE-1,i-neg);j>=0;j--)
        cur = cur*10+a[i-j]-'0';
      dig.push_back(cur);
    } trim();
  inline bool operator<(const BigInt& a)const
    {return cmp_minus(a)<0;}
  inline bool operator<=(const BigInt& a)const</pre>
    {return cmp_minus(a)<=0;}
  inline bool operator==(const BigInt& a)const
    {return cmp_minus(a)==0;}
  inline bool operator!=(const BigInt& a)const
    {return cmp_minus(a)!=0;}
  inline bool operator>(const BigInt& a)const
    {return cmp_minus(a)>0;}
  inline bool operator>=(const BigInt& a)const
    {return cmp_minus(a)>=0;}
  BigInt operator-() const {
    BigInt ret = *this;
    ret.neg ^= 1;
    return ret;
  BigInt operator+(const BigInt& a) const {
    if(neg) return -(-(*this)+(-a));
    if(a.neg) return (*this)-(-a);
    int n = max(a.len(), len());
    BigInt ret; ret.dig.resize(n);
    11d pro = 0;
    for(int i=0;i<n;i++) {</pre>
      ret.dig[i] = pro;
      if(i < a.len()) ret.dig[i] += a.dig[i];</pre>
      if(i < len()) ret.dig[i] += dig[i];</pre>
      pro = 0:
      if(ret.dig[i] >= BASE) pro = ret.dig[i]/BASE;
      ret.dig[i] -= BASE*pro;
    if(pro != 0) ret.dig.push_back(pro);
    return ret;
```

```
BigInt operator-(const BigInt& a) const {
  if(neg) return -(-(*this) - (-a));
  if(a.neg) return (*this) + (-a);
  int diff = cmp_minus(a);
  if(diff < 0) return -(a - (*this));</pre>
  if(diff == 0) return 0;
  BigInt ret; ret.dig.resize(len(), 0);
  for(int i=0;i<len();i++) {</pre>
    ret.dig[i] += dig[i];
    if(i < a.len()) ret.dig[i] -= a.dig[i];</pre>
    if(ret.dig[i] < 0){
      ret.dig[i] += BASE;
      ret.dig[i+1]--;
    }
  }
  ret.trim();
  return ret;
BigInt operator*(const BigInt& a) const {
  if(!len()||!a.len()) return 0;
  BigInt ret; ret.dig.resize(len()+a.len()+1);
  ret.neg = neg ^ a.neg;
  for(int i=0;i<len();i++)</pre>
    for(int j=0;j<a.len();j++){</pre>
      ret.dig[i+j] += dig[i] * a.dig[j];
      if(ret.dig[i+j] >= BASE) -
        lld x = ret.dig[i+j] / BASE;
        ret.dig[i+j+1] += x;
        ret.dig[i+j] -= x * BASE;
  ret.trim();
  return ret;
BigInt operator/(const BigInt& a) const {
  assert(a.len());
  if(len() < a.len()) return 0;</pre>
  BigInt ret; ret.dig.resize(len()-a.len()+1);
  ret.neg = a.neg;
  for(int i=len()-a.len();i>=0;i--){
    11d 1 = 0, r = BASE;
    while (r-1 > 1) {
      11d \ mid = (1+r)>>1;
      ret.dig[i] = mid;
      if(ret*a<=(neg?-(*this):(*this))) 1 = mid;</pre>
      else r = mid;
    ret.dig[i] = 1;
  ret.neg ^= neg; ret.trim();
  return ret;
BigInt operator%(const BigInt& a) const {
  return (*this) - (*this) / a * a;
friend BigInt abs(BigInt a){
  a.neg = 1; return a;
friend void swap(BigInt& a, BigInt& b){
  swap(a.dig, b.dig); swap(a.neg, b.neg);
friend istream& operator>>(istream& ss, BigInt& a){
  string s; ss >> s; a = s;
  return ss;
friend ostream&operator<<(ostream&o, const BigInt&a){</pre>
  if(a.len() == 0) return o << '0';
  if(a.neg) o <<</pre>
  ss << o.dig.back();
  for(int i=a.len()-2;i>=0;i--)
    o<<setw(LOG_BASE)<<setfill('0')<<a.dig[i];
  return o;
inline void print() const {
  if(len() == 0){putchar('0');return;}
  if(neg) putchar('-');
printf("%" PRINTF_ARG, dig.back());
  for(int i=len()-2;i>=0;i--)
    printf("%0" LOG_BASE_STR PRINTF_ARG, dig[i]);
#undef PRINTF_ARG
#undef LOG_BASE_STR
```

## 2.2 Dark Magic

| };

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using __gnu_pbds::pairing_heap_tag;
using __gnu_pbds::binary_heap_tag;
using __gnu_pbds::binomial_heap_tag;
using __gnu_pbds::rc_binomial_heap_tag;
using __gnu_pbds::thin_heap_tag;
template<typename T>
using pbds_heap=__gnu_pbds::prioity_queue<T, less<T>, \
                                      pairing_heap_tag>;
// a.join(b), pq.modify(pq.push(10), 87)
using __gnu_pbds::rb_tree_tag;
using __gnu_pbds::ov_tree_tag;
using __gnu_pbds::splay_tree_tag;
template<typename T>
using ordered_set = __gnu_pbds::tree<T,\</pre>
__gnu_pbds::null_type,less<T>,rb_tree_tag,\
 _gnu_pbds::tree_order_statistics_node_update>;
// find_by_order, order_of_key
template<typename A, typename B>
using hTable1=__gnu_pbds::cc_hash_table<A,B>;
template<typename A, typename B>
using hTable2=__gnu_pbds::gp_hash_table<A,B>;
```

## 2.3 Disjoint Set

```
class DJS {
private:
  vector< int > fa, sz, sv;
   vector< pair< int*, int > > opt;
   void assign( int *k, int v ) {
     opt.emplace_back( k, *k );
     *k = v;
public:
  void init( int n ) {
  fa.resize( n ); iota( fa.begin(), fa.end(), 0 );
  sz.resize( n ); fill( sz.begin(), sz.end(), 1 );
     opt.clear();
  int query(int x) {return fa[x] == x?x:query(fa[x]);}
  void merge( int a, int b ) {
     int af = query( a ), bf = query( b );
     if( af == bf ) return;
     if( sz[ af ] < sz[ bf ] ) swap( af, bf );
assign( &fa[ bf ], fa[ af ] );
assign( &sz[ af ], sz[ af ] + sz[ bf ] );</pre>
  void save() { sv.push_back( (int) opt.size() ); }
  void undo() {
     int ls = sv.back(); sv.pop_back();
     while ( ( int ) opt.size() > ls ) {
  pair< int*, int > cur = opt.back();
        *cur.first = cur.second;
        opt.pop_back();
};
```

#### 2.4 Link-Cut Tree

```
struct Node{
Node *par,*ch[2];
int xor_sum,v;
bool is_rev;
Node(int _v){
    v=xor_sum=_v;is_rev=false;
    par=ch[0]=ch[1]=nullptr;
}
inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
inline void down(){
    if(is_rev){
        if(ch[0]!=nullptr) ch[0]->set_rev();
        if(ch[1]!=nullptr) ch[1]->set_rev();
        is_rev=false;
    }
```

```
inline void up(){
    xor_sum=v;
    if(ch[0]!=nullptr){
      xor_sum^=ch[0]->xor_sum;
      ch[0]->par=this;
    if(ch[1]!=nullptr){
      xor_sum^=ch[1]->xor_sum;
      ch[1]->par=this;
  inline bool is_root(){
    return par==nullptr ||\
      (par->ch[0]!=this && par->ch[1]!=this);
  bool is_rch(){return !is_root() && par->ch[1]==this;}
} *node[maxn],*stk[maxn];
int top;
void to_child(Node* p, Node* c, bool dir) {
  p->ch[dir]=c;
  p->up();
inline void rotate(Node* node){
  Node* par=node->par;
  Node* par_par=par->par
  bool dir=node->is_rch();
  bool par_dir=par->is_rch();
  to_child(par, node->ch[!dir], dir);
  to_child(node,par,!dir);
  if(par_par!=nullptr && par_par->ch[par_dir]==par)
    to_child(par_par, node, par_dir);
  else node->par=par_par;
inline void splay(Node* node){
  Node* tmp=node;
  stk[top++]=node;
  while(!tmp->is_root()){
    tmp=tmp->par;
    stk[top++]=tmp;
  while(top) stk[--top]->down();
  for(Node *fa=node->par;
   !node->is_root();
   rotate(node), fa=node->par)
    if(!fa->is_root())
      rotate(fa->is_rch()==node->is_rch()?fa:node);
inline void access(Node* node){
  Node* last=nullptr;
  while(node!=nullptr){
    splay(node);
    to_child(node, last, true);
    last=node;
    node=node->par;
inline void change_root(Node* node){
  access(node);splay(node);node->set_rev();
inline void link(Node* x, Node* y){
  change_root(x);splay(x);x->par=y;
inline void split(Node* x, Node* y) {
  change_root(x);access(y);splay(x);
  to_child(x,nullptr,true);y->par=nullptr;
inline void change_val(Node* node,int v){
  access(node);splay(node);node->v=v;node->up();
inline int query(Node* x, Node* y)
  change_root(x);access(y);splay(y);
  return y->xor_sum;
inline Node* find_root(Node* node){
  access(node);splay(node);
  Node* last=nullptr
  while(node!=nullptr){
    node->down();last=node;node=node->ch[0];
  return last;
```

```
set<pii> dic;
inline void add_edge(int u,int v){
   if(u>v) swap(u,v);
   if(find_root(node[u])==find_root(node[v])) return;
   dic.insert(pii(u,v));
   link(node[u],node[v]);
}
inline void del_edge(int u,int v){
   if(u>v) swap(u,v);
   if(dic.find(pii(u,v))==dic.end()) return;
   dic.erase(pii(u,v));
   split(node[u],node[v]);
}
```

## 2.5 LiChao Segment Tree

```
struct Line{
 int m, k, id;
  Line() : id(-1) \{ \}
 Line( int a, int b, int c )
: m( a ), k( b ), id( c ) {}
  int at( int x ) { return m * x + k; }
class LiChao {
 private:
    int n; vector< Line > nodes;
    inline int lc( int x ) { return 2 * x + 1; }
    inline int rc( int x ) { return 2 * x + 2; }
    void insert( int 1, int r, int id, Line ln ) {
      int m = (1 + r) >> 1;
      if ( nodes[ id ].id == -1 ) {
        nodes[ id ] = ln;
        return:
      bool atLeft = nodes[ id ].at( l ) < ln.at( l );</pre>
      if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
        atLeft ^= 1; swap( nodes[ id ], ln );
      if ( r - l == 1 ) return;
      if ( atLeft ) insert( 1, m, lc( id ), ln );
      else insert( m, r, rc( id ), ln );
    int query( int 1, int r, int id, int x ) {
      int ret = 0;
      if ( nodes[ id ].id != -1 )
        ret = nodes[ id ].at( x );
      int m = (1 + r) >> 1;
      if ( r - 1 == 1 ) return ret;
      else if ( x < m )
        return max( ret, query( 1, m, lc( id ), x ) );
      else
        return max( ret, query( m, r, rc( id ), x ) );
 public:
    void build( int n_ ) {
      n = n_; nodes.clear();
      nodes.resize( n << 2, Line() );</pre>
    void insert( Line ln ) { insert( 0, n, 0, ln ); }
    int query( int x ) { return query( 0, n, 0, x ); }
} lichao;
```

## 2.6 Treap

```
namespace Treap{
 #define sz(x)((x)?((x)->size):0)
 struct node{
   int size;
   uint32_t pri;
   node *lc, *rc;
   node() : size(0), pri(rand()), lc(0), rc(0) {}
   void pull() {
     size = 1
     if ( lc ) size += lc->size;
     if ( rc ) size += rc->size;
   }
 node* merge( node* L, node* R ) {
   if ( not L or not R ) return L ? L : R;
   if ( L->pri > R->pri ) {
     L->rc = merge( L->rc, R ); L->pull();
```

```
return L:
    } else {
      R->lc = merge( L, R->lc ); R->pull();
      return R:
  void split_by_size( node*rt,int k,node*&L,node*&R ) {
    if ( not rt ) L = R = nullptr;
    else if( sz( rt->lc ) + 1 <= k ) {
      L = rt
      split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
      L->pull();
    } else {
      R = rt;
      split_by_size( rt->lc, k, L, R->lc );
      R->pull();
  #undef sz
}
```

## 2.7 SparseTable

```
template < typename T, typename Cmp_ = less< T > >
class SparseTable {
private:
  vector< vector< T > > tbl;
  vector< int > lg;
  T cv( T a, T b ) {
    return Cmp_()( a, b ) ? a : b;
public:
  void init( T arr[], int n ) {
    // 0-base
    lg.resize(n+1);
    lg[0] = -1;
    for( int i=1; i<=n; ++i) lg[i] = lg[i>>1] + 1;
    tbl.resize(lg[n] + 1);
    tbl[ 0 ].resize( n );
    copy( arr, arr + n, tbl[ 0 ].begin() );
    for ( int i = 1 ; i <= lg[ n ] ; ++ i ) {
      int len = 1 << ( i - 1 ), sz = 1 << i;</pre>
      tbl[ i ].resize( n - sz + 1 );
      for ( int j = 0 ; j <= n - sz ; ++ j
        tbl[i][j] = cv(tbl[i-1][j], tbl[i-1][j+len]);
  T query( int l, int r ) {
    // 0-base [1, r)
    int wh = lg[r - 1], len = 1 << wh;
    return cv( tbl[ wh ][ 1 ], tbl[ wh ][ r - len ] );
```

### 2.8 Linear Basis

```
struct LinearBasis {
private:
  int n, sz;
  vector< llu > B;
  inline llu two( int x ){ return ( ( llu ) 1 ) << x; }</pre>
public:
  void init( int n_ ) {
    n = n_{-}; B.clear(); B.resize( n ); sz = 0;
  void insert( llu x ) {
    // add x into B
     for ( int i = n-1; i >= 0 ; --i ) if( two(i) & x ){
       if ( B[ i ] ) x ^= B[ i ];
       else {
         B[i] = x; sz++;
         for ( int j = i - 1 ; j >= 0 ; --
         if( B[ j ] && ( two( j ) & B[ i ] ) )
B[ i ] ^= B[ j ];
for (int j = i + 1 ; j < n ; ++ j )
            if ( two( i ) & B[ j ] )
              B[ j ] ^= B[ i ];
         break:
      }
    }
```

```
inline int size() { return sz; }
bool check( llu x ) {
    // is x in span(B) ?
    for ( int i = n-1 ; i >= 0 ; --i ) if( two(i) & x )
      if( B[ i ] ) x ^= B[ i ];
      else return false;
    return true;
  llu kth_small(llu k) {
    /** 1-base would always > 0 **/
    /** should check it **/
    /* if we choose at least one element
       but size(B)(vectors in B)==N(original elements)
       then we can't get 0 */
    llu ret = 0;
    for ( int i = 0 ; i < n ; ++ i ) if( B[ i ] ) {
      if( k & 1 ) ret ^= B[ i ];
      k >>= 1;
    return ret;
} base;
```

## Graph

#### 3.1 Euler Circuit

```
bool vis[ N ]; size_t la[ K ];
void dfs( int u, vector< int >& vec ) {
 while ( la[ u ] < G[ u ].size() )</pre>
    if( vis[ G[ u ][ la[ u ] ].second ] ) {
      ++ la[ u ];
      continue;
    int v = G[ u ][ la[ u ] ].first;
    vis[ G[ u ][ la[ u ] ].second ] = true;
    ++ la[ u ]; dfs( v, vec );
    vec.push_back( v );
}
```

#### 3.2 BCC Edge

```
class BCC{
private:
  vector< int > low, dfn;
  int cnt;
  vector< bool > bridge;
  vector< vector< PII > > G;
  void dfs( int w, int f ) {
  low[ w ] = dfn[ w ] = cnt++;
    for ( auto [ u, t ] : G[ w ] ) {
  if ( u == f ) continue;
  if ( dfn[ u ] != 0 ) {
    low[ w ] = min( low[ w ], dfn[ u ] );
       }else{
         dfs( u, w );
         low[ w ] = min( low[ w ], low[ u ] );
         if ( low[ u ] > dfn[ w ] ) bridge[ t ] = true;
    }
public:
  void init( int n, int m ) {
    G.resize(n); cnt = 0;
    fill( G.begin(), G.end(), vector< PII >() );
    bridge.clear(); bridge.resize( m );
    low.clear(); low.resize( n );
dfn.clear(); dfn.resize( n );
  void add_edge( int u, int v ) {
     // should check for multiple edge
    G[ u ].emplace_back( v, cnt );
    G[ v ].emplace_back( u, cnt ++ );
  void solve(){ cnt = 1; dfs( 0, 0 ); }
  // the id will be same as insert order, 0-base
  bool is_bridge( int x ) { return bridge[ x ]; }
} bcc:
```

#### 3.3 BCC Vertex

```
class BCC{
  private:
     int n, ecnt;
     vector< vector< pair< int, int > > > G;
     vector< int > low, dfn, id;
     vector< bool > vis, ap;
     void dfs( int u, int f, int d ) {
       int child = 0;
       dfn[ u ] = low[ u ] = d; vis[ u ] = true;
       for ( auto e : G[ u ] ) if ( e.first != f ) {
         if ( vis[ e.first ] ) {
            low[ u ] = min( low[ u ], dfn[ e.first ] );
          } else {
            dfs( e.first, u, d + 1 ); child ++;
low[ u ] = min( low[ u ], low[ e.first ] );
            if ( low[ e.first ] >= d ) ap[ u ] = true;
       if ( u == f and child <= 1 ) ap[ u ] = false;</pre>
    void mark( int u, int idd ) {
   // really?????????
       if ( ap[ u ] ) return;
       for ( auto e : G[ u ] )
  if( id[ e.second ] != -1 ) {
            id[ e.second ] = idd;
            mark( e.first, idd );
  public:
    void init( int n_ ) {
       ecnt = 0, n = n_{-}
       G.clear(); G.resize( n );
       low.resize( n ); dfn.resize( n );
       ap.clear(); ap.resize( n );
       vis.clear(); vis.resize( n );
     void add_edge( int u, int v ) {
       G[ u ].emplace_back( v, ecnt );
       G[ v ].emplace_back( u, ecnt ++ );
     void solve() {
       for ( int i = 0 ; i < n ; ++ i )
  if ( not vis[ i ] ) dfs( i, i, 0 );
       id.resize( ecnt );
       fill( id.begin(), id.end(), -1 );
       for ( int i = 0 ; i < n ; ++ i )
  if ( ap[ i ] ) for ( auto e : G[ i ] )
   if( id[ e.second ] != -1 ) {</pre>
              id[ e.second ] = ecnt;
              mark( e.first, ecnt ++ );
     int get_id( int x ) { return id[ x ]; }
    int count() { return ecnt; }
bool is_ap( int u ) { return ap[ u ]; }
3.4 2-SAT (SCC)
```

```
class TwoSat{
  private:
    int n;
    vector<vector<int>> rG,G,sccs;
    vector<int> ord,idx;
    vector<bool> vis,result;
    void dfs(int u){
      vis[u]=true
      for(int v:G[u])
        if(!vis[v]) dfs(v);
      ord.push_back(u);
    void rdfs(int u){
      vis[u]=false;idx[u]=sccs.size()-1;
      sccs.back().push_back(u);
      for(int v:rG[u])
        if(vis[v])rdfs(v);
```

```
public:
    void init(int n_){
      n=n_;G.clear();G.resize(n);
      rG.clear();rG.resize(n);
      sccs.clear();ord.clear();
      idx.resize(n);result.resize(n);
    void add_edge(int u,int v){
      G[u].push_back(v);rG[v].push_back(u);
    void orr(int x,int y){
      if ((x^y)==1) return;
      add_edge(x^1,y); add_edge(y^1,x);
    bool solve(){
      vis.clear();vis.resize(n);
      for(int i=0;i<n;++i)</pre>
        if(not vis[i])dfs(i);
      reverse(ord.begin(),ord.end());
      for (int u:ord){
        if(!vis[u])continue;
        sccs.push_back(vector<int>());
        rdfs(u);
      for(int i=0;i<n;i+=2)</pre>
        if(idx[i]==idx[i+1])
          return false;
      vector<bool> c(sccs.size());
      for(size_t i=0;i<sccs.size();++i){</pre>
        for(size_t j=0;j<sccs[i].size();++j){</pre>
          result[sccs[i][j]]=c[i];
          c[idx[sccs[i][j]^1]]=!c[i];
        }
      }
      return true;
    bool get(int x){return result[x];}
    inline int get_id(int x){return idx[x];}
    inline int count(){return sccs.size();}
} sat2:
```

## 3.5 Lowbit Decomposition

```
class LowbitDecomp{
private:
  int time_, chain_, LOG_N;
  vector< vector< int > > G, fa;
  vector< int > tl, tr, chain, chain_st;
// chain_ : number of chain
  // tl, tr[ u ] : subtree interval in the seq. of u
  // chain_st[ u ] : head of the chain contains u // chian[ u ] : chain id of the chain u is on
  inline int lowbit( int x ) {
     return x & ( -x );
  void predfs( int u, int f ) {
     chain[u] = 0;
     for ( int v : G[ u ] ) {
       if ( v == f ) continue;
       predfs( v, u );
       if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )
  chain[ u ] = chain[ v ];</pre>
     if ( not chain[ u ] )
       chain[ u ] = chain_ ++;
  void dfschain( int u, int f ) {
    fa[ u ][ 0 ] = f;
for ( int i = 1 ; i < LOG_N ; ++ i )
fa[ u ][ i ] = fa[ fa[ u ][ i - 1 ] ][ i - 1 ];
     tl[u] = time_++
     if ( not chain_st[ chain[ u ] ] )
       chain_st[ chain[ u ] ] = u;
     for ( int v : G[ u ] )
       if ( v != f and chain[ v ] == chain[ u ] )
    dfschain( v, u );
for ( int v : G[ u ] )
  if ( v != f and chain[ v ] != chain[ u ] )
          dfschain( v, u );
     tr[ u ] = time_;
```

```
inline bool anc( int u, int v ) {
  return tl[ u ] <= tl[ v ] \</pre>
       and tr[ v ] <= tr[ u ];
public:
  inline int lca( int u, int v ) {
    if ( anc( u, v ) ) return u;
    for (int i = LOG_N - 1; i \ge 0; --i)
       if ( not anc( fa[ u ][ i ], v ) )
         u = fa[ u ][ i ];
    return fa[ u ][ 0 ];
  void init( int n ) {
    n ++;
    for (LOG_N = 0 ; (1 << LOG_N ) < n ; ++ LOG_N );
    fa.clear();
    fa.resize( n, vector< int >( LOG_N ) );
G.clear(); G.resize( n );
    tl.clear(); tl.resize( n );
    tr.clear(); tr.resize( n );
    chain.clear(); chain.resize( n );
    chain_st.clear(); chain_st.resize( n );
  void add_edge( int u , int v ) {
    // 1-base
    G[ u ].push_back( v );
    G[ v ].push_back( u );
  void decompose(){
    chain_ = 1;
    predfs( 1, 1 );
    time_{-} = 0;
    dfschain( 1, 1 );
  PII get_inter( int u ) { return {tl[ u ], tr[ u ]}; }
  vector< PII > get_path( int u , int v ){
  vector< PII > res;
    int g = lca( u, v );
    while ( chain[ u ] != chain[ g ] ) {
  int s = chain_st[ chain[ u ] ];
       res.emplace_back( tl[ s ], tl[ u ] + 1 );
       u = fa[ s ][ 0 ];
    res.emplace_back(tl[g], tl[u] + 1);
    while ( chain[ v ] != chain[ g ] ) {
  int s = chain_st[ chain[ v ] ];
       res.emplace_back( tl[ s ], tl[ v ] + 1 );
       v = fa[ s ][ 0 ];
    res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
     return res;
    /* res : list of intervals from u to v
     * ( note only nodes work, not edge )
      * vector< PII >& path = tree.get_path( u , v )
      * for( auto [ 1, r ] : path ) {
         0-base [ 1, r )
      * }
} tree;
3.6 MaxClique
```

```
for( size_t u = G[ i ]._Find_first()
          u < n ; u = G[ i ]._Find_next( u ) )
            -- deg[ u ];
  void BK( bits R, bits P, bits X ) {
    if (R.count()+P.count() <= ans.count()) return;</pre>
    if ( not P.count() and not X.count() ) {
      if ( R.count() > ans.count() ) ans = R;
      return:
    /* greedily chosse max degree as pivot
    bits cur = P | X; size_t pivot = 0, sz = 0;
    for ( size_t u = cur._Find_first() ;
      u < n ; u = cur._Find_next( u )</pre>
        if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
    cur = P & ( \sim G[ pivot ] );
    ^{*}/ // or simply choose first
    bits cur = P & (~G[ ( P | X )._Find_first() ]);
    for ( size_t u = cur._Find_first() ;
      u < n ; u = cur._Find_next( u ) ) {
      if ( R[ u ] ) continue;
      R[u] = 1;
      BK( R, P & G[ u ], X & G[ u ]);
      R[u] = P[u] = 0, X[u] = 1;
public:
  void init( size_t n_ ) {
    n = n_{-};
    for ( size_t i = 0 ; i < n ; ++ i )</pre>
      G[ i ].reset();
    ans.reset();
  void add_edges( int u, bits S ) { G[ u ] = S; }
  void add_edge( int u, int v ) {
    G[u][v] = G[v][u] = 1;
  int solve() {
    sort_by_degree(); // or simply iota( deo... )
    for ( size_t i = 0 ; i < n ; ++ i )</pre>
      deg[ i ] = G[ i ].count();
    bits pob, nob = 0; pob.set()
    for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
    for ( size_t i = 0 ; i < n ; ++ i ) {
      size_t v = deo[ i ];
      bits tmp; tmp[v] = 1;
      BK( tmp, pob & G[ v ], nob & G[ v ]);
      pob[v] = 0, nob[v] = 1;
    return static_cast< int >( ans.count() );
};
      Virtural Tree
inline bool cmp(const int &i, const int &j) {
  return dfn[i] < dfn[j];</pre>
void build(int vectrices[], int k) {
  static int stk[MAX_N];
  sort(vectrices, vectrices + k, cmp);
```

## 3.8 Tree Hashing

```
uint64_t hsah( int u, int f ) {
    uint64_t r = 127;
    for ( int v : G[ u ] ) {
        if ( v == f ) continue;
        uint64_t hh = hsah( v, u );
        r = r + ( hh * hh ) % mod;
    }
    return r;
}
```

## 3.9 Minimum Mean Cycle

```
/* minimum mean cycle O(VE) */
struct MMC{
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
#define V 1021
#define inf 1e9
   struct Edge { int v,u; double c; };
   int n, m, prv[V][V], prve[V][V], vst[V];
   Edge e[E];
   vector<int> edgeID, cycle, rho;
   double d[V][V];
   void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
   void add_edge( int vi , int ui , double ci )
   { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
     for(int i=0; i<n; i++) d[0][i]=0;</pre>
     for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);</pre>
       for(int j=0; j<m; j++)</pre>
         int v = e[j].v, u = e[j].u;
          if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
            d[i+1][u] = d[i][v]+e[j].c;
            prv[i+1][u] = v;
            prve[i+1][u] = j;
       }
    }
   }
   double solve(){
     // returns inf if no cycle, mmc otherwise
     double mmc=inf;
     int st = -1;
     bellman_ford();
     for(int i=0; i<n; i++) {</pre>
       double avg=-inf;
       for(int k=0; k<n; k++) {</pre>
          if(d[n][i]<inf-eps)</pre>
            avg=max(avg,(d[n][i]-d[k][i])/(n-k));
         else avg=max(avg,inf);
       if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
     FZ(vst);edgeID.clear();cycle.clear();rho.clear();
     for (int i=n; !vst[st]; st=prv[i--][st]) {
       vst[st]++;
       edgeID.PB(prve[i][st]);
       rho.PB(st);
     while (vst[st] != 2) {
       int v = rho.back(); rho.pop_back();
       cycle.PB(v);
       vst[v]++;
     reverse(ALL(edgeID));
     edgeID.resize(SZ(cycle));
     return mmc;
} mmc;
```

## 3.10 Mo's Algorithm on Tree

```
int n, q, nxt[ N ], to[ N ], hd[ N ];
struct Que{
  int u, v, id;
} que[ N ];
void init() {
```

```
cin >> n >> q;
for ( int i = 1 ; i < n ; ++ i ) {</pre>
    int u, v; cin >> u >> v;
    nxt[ i << 1 | 0 ] = hd[ u ];</pre>
    to[i << 1 | 0] = v;
    hd[u] = i << 1 | 0;
    nxt[i << 1 | 1] = hd[v];
    to[i << 1 | 1] = u;
    hd[v] = i << 1 | 1;
  for ( int i = 0 ; i < q ; ++ i ) {
    cin >> que[ i ].u >> que[ i ].v; que[ i ].id = i;
  }
int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
void dfs( int u, int f ) {
  dfn[ u ] = dfn_++; int saved_rbp = stk_;
for ( int v_ = hd[ u ] ; v_ ; v_ = nxt[ v_ ] ) {
    if ( to[ v_ ] == f ) continue;
    dfs( to[v_{-}], u);
    if ( stk_ - saved_rbp < SQRT_N ) continue;</pre>
    for ( ++ block_ ; stk_ != saved_rbp ; )
       block_id[ stk[ -- stk_ ] ] = block_;
  stk[ stk_+ ++ ] = u;
bool inPath[ N ];
void Diff( int u ) {
  if ( inPath[ u ] ^= 1 )
    // remove this edge
  else
    // add this edge
void traverse( int& origin_u, int u ) {
  for ( int g = lca( origin_u, u ) ;
    origin_u != g ; origin_u = parent_of[ origin_u ] )
      Diff( origin_u );
  for (int v = u; v != origin_u; v = parent_of[v])
    Diff( v );
  origin_u = u;
void solve() {
  while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
  sort( que, que + q, [](const Que& x, const Que& y) {
    return tie( block_id[ x.u ], dfn[ x.v ] )
             < tie( block_id[ y.u ], dfn[ y.v ] );
  } );
  int U = 1, V = 1;
for ( int i = 0 ; i < q ; ++ i ) {
    pass( U, que[ i ].u );
pass( V, que[ i ].v );
    // we could get our answer of que[ i ].id
}
/*
Method 2:
dfs u:
 push u
  iterate subtree
 push u
Let P = LCA(u, v), and St(u) <= St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
```

## 3.11 Minimum Steiner Tree

```
// Minimum Steiner Tree
// O(V 3^T + V^2 2^T)
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
  int n , dst[V][V] , dp[1 << T][V] , tdst[V];
  void init( int _n ){
    n = _n;
    for( int i = 0 ; i < n ; i ++ ){
        for( int j = 0 ; j < n ; j ++ )
            dst[ i ][ j ] = INF;
        dst[ i ][ i ] = 0;</pre>
```

```
}
  void add_edge( int ui , int vi , int wi ){
     dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
     dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
  void shortest_path(){
     for( int k = 0 ; k < n ; k ++ )
       for( int i = 0 ; i < n ; i ++ )</pre>
         for( int j = 0 ; j < n ; j ++ )
  dst[ i ][ j ] = min( dst[ i ][ j ],</pre>
                   dst[ i ][ k ] + dst[ k ][ j ] );
  int solve( const vector<int>& ter ){
     int t = (int)ter.size();
     for( int i = 0 ; i < ( 1 << t ) ; i ++ )
       for( int j = 0 ; j < n ; j ++ )
  dp[ i ][ j ] = INF;</pre>
     for( int i = 0 ; i < n ; i ++ )
       dp[0][i] = 0;
     for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
       if( msk == ( msk & (-msk) ) ){
         int who = __lg( msk );
for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];</pre>
         continue;
       for( int i = 0 ; i < n ; i ++ )</pre>
          for( int submsk = ( msk - 1 ) & msk ; submsk ;
    submsk = ( submsk - 1 ) & msk )
               dp[ msk ][ i ] = min( dp[ msk ][ i ],
                                  dp[ submsk ][ i ] +
                                  dp[ msk ^ submsk ][ i ] );
       for( int i = 0 ; i < n ; i ++ ){
          tdst[ i ] = INF;
          for( int j = 0 ; j < n ; j ++ )
  tdst[ i ] = min( tdst[ i ],</pre>
                          dp[ msk ][ j ] + dst[ j ][ i ] );
       for( int i = 0 ; i < n ; i ++ )</pre>
          dp[ msk ][ i ] = tdst[ i ];
     int ans = INF;
     for( int i = 0 ; i < n ; i ++ )</pre>
       ans = min(ans, dp[(1 << t) - 1][i]);
     return ans:
} solver;
```

## 3.12 Directed Minimum Spanning Tree

```
template <typename T> struct DMST {
  T g[maxn][maxn], fw[maxn];
  int n, fr[maxn];
  bool vis[maxn], inc[maxn];
  void clear() {
    for(int i = 0; i < maxn; ++i) {</pre>
      for(int j = 0; j < maxn; ++j) g[i][j] = inf;
      vis[i] = inc[i] = false;
  void addEdge(int u,int v,T w){g[u][v]=min(g[u][v],w)
      ;}
 T operator()(int root, int _n) {
  n = _n; T ans = 0;
    if (dfs(root) != n) return -1;
    while (true) {
      for(int i = 1;i <= n;++i) fw[i] = inf, fr[i] = i;
      for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
        for (int j = 1; j <= n; ++j) {</pre>
          if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
            fw[i] = g[j][i]; fr[i] = j;
          }
        }
      int x = -1;
      for(int i = 1;i <= n;++i)if(i != root && !inc[i])</pre>
        int j = i, c = 0;
        while(j!=root && fr[j]!=i && c<=n) ++c, j=fr[j</pre>
```

```
if (j == root || c > n) continue;
         else { x = i; break; }
      if (!~x) {
         for (int i = 1; i <= n; ++i)
           if (i != root && !inc[i]) ans += fw[i];
         return ans;
      int y = x;
      for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
         ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true
       } while (y != x);
      inc[x] = false;
       for (int k = 1; k <= n; ++k) if (vis[k]) {
         for (int j = 1; j <= n; ++j) if (!vis[j]) {
  if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
           if (g[j][k] < \inf \&\& g[j][k]-fw[k] < g[j][x])
             g[j][x] = g[j][k] - fw[k];
      }
    return ans;
  int dfs(int now) {
    int r = 1; vis[now] = true;
    for (int i = 1; i <= n; ++i)</pre>
       if (g[now][i] < inf && !vis[i]) r += dfs(i);</pre>
     return r;
  }
};
```

## 3.13 Dominator Tree

```
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
  // vertices are numbered from 0 to n-1
  fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
  fill(fa, fa + n, -1); fill(val, val + n, -1);
  fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
  fill(dom, dom + n, -1); tk = 0; for (int i = 0; i < n; ++i) {
    g[i].clear(); r[i].clear(); rdom[i].clear();
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
  rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
  for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
    r[dfn[u]].push_back(dfn[x]);
  }
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
  if (fa[x] == x) return c ? -1 : x;
  int p = find(fa[x], 1);
  if (p == -1) return c ? fa[x] : val[x];
  if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
  fa[x] = p;
  return c ? p : val[x];
vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
// p[i] = -2 if i is unreachable from s
  dfs(s);
  for (int i = tk - 1; i >= 0; --i) {
    for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)])
    if (i) rdom[sdom[i]].push_back(i);
    for (int &u : rdom[i]) {
      int p = find(u);
      if (sdom[p] == i) dom[u] = i;
      else dom[u] = p;
    if (i) merge(i, rp[i]);
```

```
}
vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];
return p;
}}</pre>
```

## 4 Matching & Flow

#### 4.1 Kuhn Munkres

```
class KM {
private:
  static constexpr lld INF = 1LL << 60;</pre>
  vector<lld> hl,hr,slk;
  vector<int> f1,fr,pre,qu;
  vector<vector<lld>> w;
  vector<bool> v1,vr;
  int n, ql, qr;
  bool check(int x) {
    if (vl[x] = true, fl[x] != -1)
      return vr[qu[qr++] = f1[x]] = true;
    while (x != -1) swap(x, fr[fl[x] = pre[x]]);
    return false:
  void bfs(int s) {
    fill(slk.begin(), slk.end(), INF);
    fill(vl.begin(), vl.end(), false);
    fill(vr.begin(), vr.end(), false);
    ql = qr = 0;
    qu[qr++] = s;
    vr[s] = true;
    while (true) {
      11d d;
      while (ql < qr) {</pre>
        for (int x = 0, y = qu[ql++]; x < n; ++x)
           if(!v1[x]\&\&s1k[x]>=(d=h1[x]+hr[y]-w[x][y]))\{
             if (pre[x] = y, d) slk[x] = d;
             else if (!check(x)) return;
        }
      d = INF:
      for (int x = 0; x < n; ++x)
        if (!v1[x] \&\& d > s1k[x]) d = s1k[x];
       for (int x = 0; x < n; ++x) {
        if (v1[x]) h1[x] += d;
         else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
         if (!v1[x] && !slk[x] && !check(x)) return;
    }
public:
  void init( int n_ ) {
    n = n_{;} qu.resize(n);
    fl.clear(); fl.resize(n, -1);
fr.clear(); fr.resize(n, -1);
    hr.clear(); hr.resize(n); hl.resize(n);
    w.clear(); w.resize(n, vector<lld>(n));
    slk.resize(n); pre.resize(n);
    vl.resize(n); vr.resize(n);
  void set_edge( int u, int v, lld x ) {w[u][v] = x;}
  1ld solve() {
    for (int i = 0; i < n; ++i)
      hl[i] = *max_element(w[i].begin(), w[i].end());
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    11d res = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];
    return res;
} km;
```

## 4.2 Bipartite Matching

```
class BipartiteMatching{
private:
```

```
vector<int> X[N], Y[N];
int fX[N], fY[N], n;
  bitset<N> walked;
  bool dfs(int x){
     for(auto i:X[x]){
       if(walked[i])continue;
       walked[i]=1;
       if(fY[i]==-1||dfs(fY[i])){
         fY[i]=x;fX[x]=i;
         return 1;
     }
     return 0;
public:
  void init(int _n){
    n=_n; walked.reset();
     for(int i=0;i<n;i++){</pre>
       X[i].clear();Y[i].clear();
       fX[i]=fY[i]=-1;
     }
  void add_edge(int x, int y){
  X[x].push_back(y); Y[y].push_back(y);
  int solve(){
     int cnt = 0;
     for(int i=0;i<n;i++){</pre>
       walked.reset();
       if(dfs(i)) cnt++;
     // return how many pair matched
     return cnt:
  }
};
```

#### 4.3 General Graph Matching

```
const int N = 514, E = (2e5) * 2;
struct Graph{
  int to[E],bro[E],head[N],e;
  int lnk[N], vis[N], stp, n;
  void init( int _n ){
    stp = 0; e = 1; n = _n;
for( int i = 0 ; i <= n ; i ++ )
      head[i] = lnk[i] = vis[i] = 0;
  void add_edge(int u,int v){
    // 1-base
    to[e]=v,bro[e]=head[u],head[u]=e++;
    to[e]=u, bro[e]=head[v], head[v]=e++;
  bool dfs(int x){
    vis[x]=stp;
    for(int i=head[x];i;i=bro[i]){
      int v=to[i]:
      if(!lnk[v]){
         lnk[x]=v, lnk[v]=x;
         return true
      }else if(vis[lnk[v]]<stp){</pre>
         int w=lnk(v);
        lnk[x]=v, lnk[v]=x, lnk[w]=0;
         if(dfs(w)) return true;
        lnk[w]=v, lnk[v]=w, lnk[x]=0;
      }
    return false;
  int solve(){
    int ans = 0;
    for(int i=1;i<=n;i++)</pre>
      if(not lnk[i]){
        stp++; ans += dfs(i);
    return ans;
} graph;
```

#### Minimum Weight Matching (Clique 4.4 version)

```
struct Graph {
  // 0-base (Perfect Match)
  int n, edge[MXN][MXN];
  int match[MXN], dis[MXN], onstk[MXN];
  vector<int> stk;
  void init(int _n) {
    n = _n;
for (int i=0; i<n; i++)</pre>
       for (int j=0; j<n; j++)</pre>
         edge[i][j] = 0;
  void set_edge(int u, int v, int w) {
     edge[u][v] = edge[v][u] = w;
  bool SPFA(int u){
    if (onstk[u]) return true;
    stk.PB(u);
     onstk[u] = 1;
     for (int v=0; v<n; v++){
       if (u != v && match[u] != v && !onstk[v]){
         int m = match[v];
         if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
           dis[m] = dis[u] - edge[v][m] + edge[u][v];
           onstk[v] = 1;
           stk.PB(v);
           if (SPFA(m)) return true;
           stk.pop_back();
           onstk[v] = 0;
         }
      }
    onstk[u] = 0;
    stk.pop_back();
     return false;
  int solve() {
     // find a match
    for (int i=0; i<n; i+=2){
  match[i] = i+1;</pre>
      match[i+1] = i;
     while (true){
       int found = 0;
       for (int i=0; i<n; i++)
         dis[i] = onstk[i] = 0;
       for (int i=0; i<n; i++){</pre>
         stk.clear()
         if (!onstk[i] && SPFA(i)){
           found = 1:
           while (SZ(stk)>=2){
             int u = stk.back(); stk.pop_back();
             int v = stk.back(); stk.pop_back();
             match[u] = v;
             match[v] = u;
         }
       if (!found) break;
     int ret = 0;
     for (int i=0; i<n; i++)
       ret += edge[i][match[i]];
     return ret>>1;
} graph;
```

#### 4.5 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x,y,l,u), connect x o y with capacity
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S \to v$  with capacity in(v), otherwise, connect v o T with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from  $\boldsymbol{s}$  to  $\boldsymbol{t}$  is the answer.

```
To minimize, let f be the maximum flow from S to
T. Connect t \to s with capacity \infty and let the flow from S to T be f'. If f+f' \neq \sum_{v \in V, in(v)>0} in(v), there's no solution. Otherwise, f' is the answer.
```

- 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching  ${\cal M}$  on bipartite graph (X, Y)
  - 1. Redirect every edge:  $y \to x$  if  $(x,y) \in M$ ,  $x \to y$  otherwise.
  - 2. DFS from unmatched vertices in X. 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source  ${\it S}$  and sink  ${\it T}$
  - 2. For each edge (x,y,c), connect  $x \to y$  with (cost,cap) =(c,1) if c>0, otherwise connect  $y \to x$  with (cost, cap) =
  - 3. For each edge with c < 0, sum these cost as K, then
  - increase d(y) by 1, decrease d(x) by 1 4. For each vertex v with d(v)>0, connect  $S\to v$  with
  - (cost, cap) = (0, d(v)) 5. For each vertex v with d(v) < 0, connect  $v \to T$  with (cost, cap) = (0, -d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C + K
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer
  - 2. Construct a max flow model, let K be the sum of all weights
  - 3. Connect source  $s \to v$ ,  $v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$ with capacity  $\boldsymbol{w}$
  - 5. For  $v \in G$ , connect it with sink  $v \to t$  with capacity  $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
  - 6. T is a valid answer if the maximum flow f < K |V|
- · Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect u' o v'
  - with weight w(u,v). 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 3. Find the minimum weight perfect matching on  $G^\prime$ .
- Project selection problem
  - 1. If  $p_v>0$ , create edge (s,v) with capacity  $p_v$ ; otherwise,
  - create edge (v,t) with capacity  $-p_v$ . 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v. 3. The mincut is equivalent to the maximum profit of a
  - subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity  $c_x$  and create edge (s,y)
- with capacity  $c_y$ . 2. Create edge (x,y) with capacity  $c_{xy}$ . 3. Create edge (x,y) and edge (x',y') with capacity  $c_{xyx'y'}$ .

#### 4.6 Dinic

```
class Dinic{
private:
  using CapT = int64_t;
  struct Edge{
    int to, rev;
    CapT cap;
  int n, st, ed;
  vector<vector<Edge>> G;
  vector<int> lv, idx;
  bool BFS(){
    fill(lv.begin(), lv.end(), -1);
    queue<int> bfs;
    bfs.push(st);
    lv[st] = 0;
    while(!bfs.empty()){
      int u = bfs.front(); bfs.pop();
      for(auto e: G[u]){
        if(e.cap <= 0 or lv[e.to]!=-1) continue;
lv[e.to] = lv[u] + 1;</pre>
        bfs.push(e.to);
```

```
return (lv[ed]!=-1);
  CapT DFS(int u, CapT f){
     if(u == ed) return f;
     CapT ret = 0;
     for(int& i = idx[u]; i < (int)G[u].size(); ++i){</pre>
       auto& e = G[u][i];
       if(e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
       CapT nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
       if(f == 0) return ret;
     if(ret == 0) lv[u] = -1;
     return ret;
public:
  void init(int n_, int st_, int ed_){
     n = n_{,} st = st_{,} ed = ed_{,}
     G.resize(n); lv.resize(n);
     fill(G.begin(), G.end(), vector<Edge>());
  void add_edge(int u, int v, CapT c){
  G[u].push_back({v, (int)G[v].size(), c});
  G[v].push_back({u, ((int)G[u].size())-1, 0});
  CapT max_flow(){
     CapT ret = 0;
     while(BFS()){
       idx.assign(n, 0);
       CapT f = DFS(st, numeric_limits<CapT>::max());
       ret += f:
       if(f == 0) break;
     return ret;
} flow;
```

#### 4.7 Minimum Cost Maximum Flow

```
class MiniCostMaxiFlow{
  using CapT = int;
  using WeiT = int64_t;
  using PCW = pair<CapT,WeiT>;
  static constexpr CapT INF_CAP = 1 << 30;
static constexpr WeiT INF_WEI = 1LL<<60;</pre>
private:
  struct Edge{
    int to, back;
    WeiT wei;
    CapT cap;
    Edge() {}
    Edge(int a,int b,WeiT c,CapT d):
      to(a), back(b), wei(c), cap(d)
    {}
  }:
  int ori, edd;
  vector<vector<Edge>> G;
  vector<int> fa, wh;
  vector<bool> inq;
  vector<WeiT> dis;
  PCW SPFA(){
    fill(inq.begin(),inq.end(),false);
    fill(dis.begin(), dis.end(), INF_WEI);
    queue<int> qq; qq.push(ori);
    dis[ori]=0;
    while(!qq.empty()){
       int u=qq.front();qq.pop();
       inq[u] = 0;
       for(int i=0;i<SZ(G[u]);++i){</pre>
         Edge e=G[u][i];
         int v=e.to;
         WeiT d=e.wei;
         if(e.cap <= 0 | |dis[v] <= dis[u] + d)
           continue
         dis[v]=dis[u]+d;
         fa[v]=u, wh[v]=i;
         if(ing[v]) continue;
         qq.push(v);
         inq[v]=1;
```

```
if(dis[edd]==INF_WEI)
      return {-1,-1};
    CapT mw=INF_CAP;
    for(int i=edd;i!=ori;i=fa[i])
      mw=min(mw,G[fa[i]][wh[i]].cap);
    for (int i=edd;i!=ori;i=fa[i]){
      auto &eg=G[fa[i]][wh[i]];
      eg.cap-=mw;
      G[eg.to][eg.back].cap+=mw;
    return {mw,dis[edd]};
public:
 void init(int a,int b,int n){
    ori=a,edd=b;
    G.clear();G.resize(n);
    fa.resize(n);wh.resize(n);
    inq.resize(n); dis.resize(n);
  void add_edge(int st,int ed,WeiT w,CapT c){
    G[st].emplace_back(ed,SZ(G[ed]),w,c);
    G[ed].emplace_back(st,SZ(G[st])-1,-w,0);
  PCW solve(){
    /* might modify to
    cc += ret.first * ret.second
    ww += ret.first * ret.second
    CapT cc=0; WeiT ww=0;
    while(true){
      PCW ret=SPFA();
      if(ret.first==-1) break;
      cc+=ret.first;
      ww+=ret.second;
    return {cc,ww};
} mcmf:
```

## 4.8 Global Min-Cut

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
    w[x][y] += c;
    w[y][x] += c;
pair<int, int> phase(int n) {
    memset(v, false, sizeof(v));
    memset(g, 0, sizeof(g));
int s = -1, t = -1;
    while (true) {
        int c = -1;
        for (int i = 0; i < n; ++i) {</pre>
             if (del[i] || v[i]) continue;
             if (c == -1 \mid | g[i] > g[c]) c = i;
        if (c == -1) break;
        v[c] = true;
        s = t, t = c;
         for (int i = 0; i < n; ++i) {
             if (del[i] || v[i]) continue;
             g[i] += w[c][i];
    return make_pair(s, t);
}
int mincut(int n) {
    int cut = 1e9;
    memset(del, false, sizeof(del));
for (int i = 0; i < n - 1; ++i) {</pre>
         int s, t; tie(s, t) = phase(n);
        del[t] = true;
        cut = min(cut, g[t]);
        for (int j = 0; j < n; ++j) {
```

```
w[s][j] += w[t][j];
w[j][s] += w[j][t];
}

return cut;
}
```

## 5 Math

#### 5.1 Prime Table

```
\begin{array}{c} 1002939109, 1020288887, 1028798297, 1038684299, \\ 1041211027, 1051762951, 1058585963, 1063020809, \\ 1147930723, 1172520109, 1183835981, 1187659051, \\ 1241251303, 1247184097, 1255940849, 1272759031, \\ 1287027493, 1288511629, 1294632499, 1312650799, \\ 1868732623, 1884198443, 1884616807, 1885059541, \\ 1909942399, 1914471137, 1923951707, 1925453197, \\ 1979612177, 1980446837, 1989761941, 2007826547, \\ 2008033571, 2011186739, 2039465081, 2039728567, \\ 2093735719, 2116097521, 2123852629, 2140170259, \\ 3148478261, 3153064147, 3176351071, 3187523093, \\ 3196772239, 3201312913, 3203063977, 3204840059, \\ 3210224309, 3213032591, 3217689851, 3218469083, \\ 3219857533, 3231880427, 3235951699, 3273767923, \\ 3276188869, 3277183181, 3282463507, 3285553889, \\ 3319309027, 3327005333, 3327574903, 3341387953, \\ 3373293941, 3380077549, 3380892997, 3381118801 \\ \end{array}
```

# 5.2 $\lfloor \frac{n}{i} \rfloor$ Enumeration

```
T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor
```

## 5.3 ax+by=qcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b==(x/y)*a;
}
```

#### 5.4 Pollard Rho

```
// does not work when n is prime
// return any non-trivial factor
llu pollard_rho(llu n){
    static auto f=[](llu x,llu k,llu m){
        return add(k,mul(x,x,m),m);
    };
    if (!(n&1)) return 2;
    mt19937 rnd(120821011);
    while(true){
        llu y=2,yy=y,x=rnd()%n,t=1;
        for(llu sz=2;t==1;sz<<=1) {
            for(llu i=0;i<sz;++i){
                if(t!=1)break;
                yy=f(yy,x,n);
                t=gcd(yy>y?yy-y:y-yy,n);
            }
            y=yy;
        }
        if(t!=1&&t!=n) return t;
    }
}
```

## 5.5 Pi Count (Linear Sieve)

```
static constexpr int N = 1000000 + 5;
lld pi[N];
vector<int> primes;
bool sieved[N];
lld cube_root(lld x){
    lld s=cbrt(x-static_cast<long double>(0.1));
    while(s*s*s <= x) ++s;
    return s-1;
}
lld square_root(lld x){
    lld s=sqrt(x-static_cast<long double>(0.1));
    while(s*s <= x) ++s;
    return s-1;
}
void init(){
    primes.reserve(N);</pre>
```

```
primes.push_back(1);
  for(int i=2;i<N;i++) {</pre>
    if(!sieved[i]) primes.push_back(i);
    pi[i] = !sieved[i] + pi[i-1];
    for(int p: primes) if(p > 1) {
      if(p * i >= N) break;
      sieved[p * i] = true;
      if(p % i == 0) break;
 }
lld phi(lld m, lld n) {
  static constexpr int MM = 80000, NN = 500;
  static lld val[MM][NN];
  \label{lem:mass}  \textbf{if}(m < MM\&\&n < NN\&\&val[m][n]) \\ return \\ val[m][n]-1; 
  if(n == 0) return m;
 if(primes[n] >= m) return 1;
lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
  if(m < MM\&n < NN) val[m][n] = ret+1;
  return ret;
1ld pi_count(1ld);
11d P2(11d m, 11d n) {
  11d sm = square_root(m), ret = 0;
  for(lld i = n+1;primes[i]<=sm;i++)</pre>
    ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
  return ret;
11d pi_count(11d m) {
  if(m < N) return pi[m];</pre>
  11d n = pi_count(cube_root(m));
  return phi(m, n) + n - 1 - P2(m, n);
```

## 5.6 Range Sieve

```
const int MAX_SQRT_B = 50000;
const int MAX_L = 200000 + 5;
bool is_prime_small[MAX_SQRT_B];
bool is_prime[MAX_L];
void sieve(lld l, lld r){
  // [1, r)
  for(lld i=2;i*i<r;i++) is_prime_small[i] = true;</pre>
  for(lld i=1;i<r;i++) is_prime[i-1] = true;</pre>
  if(l==1) is_prime[0] = false;
  for(lld i=2;i*i<r;i++){</pre>
    if(!is_prime_small[i]) continue;
    for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;</pre>
    for(lld j=std::max(2LL, (1+i-1)/i)*i;j<r;j+=i)
        is_prime[j-1]=false;
  }
}
```

## 5.7 Miller Rabin

```
bool isprime(llu x){
 static llu magic[]={2,325,9375,28178,\
                    450775,9780504,1795265022};
  static auto witn=[](llu a,llu u,llu n,int t){
    a = mpow(a,u,n);
    if (!a)return 0;
    while(t--){
      1lu a2=mul(a,a,n);
      if(a2==1 && a!=1 && a!=n-1)
        return 1;
      a = a2;
    }
    return a!=1;
  if(x<2)return 0;</pre>
  if(!(x&1))return x==2;
 llu x1=x-1;int t=0;
 while(!(x1&1))x1>>=1,t++;
  for(llu m:magic)if(witn(m,x1,x,t))return 0;
  return 1;
```

## 5.8 Inverse Element

```
// x's inverse mod k
long long GetInv(long long x, long long k){
   // k is prime: euler_(k)=k-1
   return qPow(x, euler_phi(k)-1);
}
// if you need [1, x] (most use: [1, k-1]
void solve(int x, long long k){
   inv[1] = 1;
   for(int i=2;i<x;i++)
     inv[i] = ((long long)(k - k/i) * inv[k % i]) % k;
}</pre>
```

#### 5.9 Euler Phi Function

```
extended euler:
   a^b mod p
   if gcd(a, p)==1: a^(b%phi(p))
elif b < phi(p): a^b mod p</pre>
   else a^(b%phi(p) + phi(p))
lld euler_phi(int x){
  11d r=1;
  for(int i=2;i*i<=x;++i){</pre>
    if(x\%i==0){
       x/=i; r*=(i-1);
       while(x%i==0) {
        x/=i; r*=i;
  if(x>1) r*=x-1;
  return r;
vector<int> primes;
bool notprime[N];
11d phi[N];
void euler_sieve(int n){
  for(int i=2;i<n;i++){</pre>
    if(!notprime[i]){
       primes.push_back(i); phi[i] = i-1;
    for(auto j: primes){
       if(i*j >= n) break;
       notprime[i*j] = true;
       phi[i*j] = phi[i] * phi[j];
       if(i % j == 0){
         phi[i*j] = phi[i] * j;
         break:
    }
  }
}
```

## 5.10 Gauss Elimination

```
void gauss(vector<vector<double>> &d) {
   int n = d.size(), m = d[0].size();
   for (int i = 0; i < m; ++i) {
      int p = -1;
      for (int j = i; j < n; ++j) {
        if (fabs(d[j][i]) < eps) continue;
        if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
      }
      if (p == -1) continue;
      for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
      for (int j = 0; j < n; ++j) {
        if (i == j) continue;
        double z = d[j][i] / d[i][i];
        for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
      }
    }
}</pre>
```

#### 5.11 Fast Fourier Transform

```
/*
  polynomial multiply:
```

```
DFT(a, len); DFT(b, len);
for(int i=0;i<len;i++) c[i] = a[i]*b[i];</pre>
   iDFT(c, len);
   (len must be 2^k and = 2^k (max(a, b)))
   Hand written Cplx would be 2x faster
*/
Cplx omega[2][N];
void init_omega(int n) {
 static constexpr llf PI=acos(-1);
  const llf arg=(PI+PI)/n;
  for(int i=0;i<n;++i)</pre>
    omega[0][i]={cos(arg*i), sin(arg*i)};
  for(int i=0;i<n;++i)</pre>
    omega[1][i]=conj(omega[0][i]);
void tran(Cplx arr[],int n,Cplx omg[]) {
  for(int i=0, j=0;i<n;++i){</pre>
    if(i>j)swap(arr[i],arr[j]);
    for(int l=n>>1;(j^=1)<1;l>>=1);
  for (int l=2;l<=n;l<<=1){</pre>
    int m=1>>1;
    for(auto p=arr;p!=arr+n;p+=1){
      for(int i=0;i<m;++i){</pre>
        Cplx t=omg[n/1*i]*p[m+i];
        p[m+i]=p[i]-t; p[i]+=t;
    }
 }
void DFT(Cplx arr[],int n){tran(arr,n,omega[0]);}
void iDFT(Cplx arr[],int n){
  tran(arr,n,omega[1]);
  for(int i=0;i<n;++i) arr[i]/=n;</pre>
```

## 5.12 High Speed Linear Recurrence

```
#define mod 998244353
const int N=1000010;
int n,k,m,f[N],h[N],a[N],b[N],ib[N];
int pw(int x,int y){
  int re=1:
  if(y<0)y+=mod-1;
  while(y){
    if(y&1)re=(11)re*x%mod;
   y >> = 1; x = (11)x*x%mod;
  return re;
void inc(int&x,int y){x+=y;if(x>=mod)x-=mod;}
namespace poly{
 const int G=3;
  int rev[N],L;
  void ntt(int*A,int len,int f){
    for(L=0;(1<<L)<len;++L);
    for(int i=0;i<len;++i){</pre>
      rev[i]=(rev[i>>1]>>1)|((i&1)<<(L-1));
      if(i<rev[i])swap(A[i],A[rev[i]]);</pre>
    for(int i=1;i<len;i<<=1){</pre>
      int wn=pw(G, f*(mod-1)/(i<<1));</pre>
      for(int j=0;j<len;j+=i<<1){</pre>
        int w=1;
        for(int k=0;k<i;++k,w=(11)w*wn%mod){</pre>
           int x=A[j+k], y=(11)w*A[j+k+i]%mod;
          A[j+k]=(x+y) \mod A[j+k+i]=(x-y+mod) \mod ;
      }
    if(!~f){
      int iv=pw(len,mod-2);
      for(int i=0;i<len;++i)A[i]=(11)A[i]*iv%mod;</pre>
  void cls(int*A,int l,int r){
    for(int i=1;i<r;++i)A[i]=0;}</pre>
  void cpy(int*A,int*B,int 1){
    for(int i=0;i<1;++i)A[i]=B[i];}</pre>
  void inv(int*A, int*B, int 1){
    if(l==1){B[0]=pw(A[0],mod-2);return;}
```

```
static int t[N];
    int len=1<<1;
    inv(A,B,l>>1);
    cpy(t, A, 1); cls(t, 1, len);
    ntt(t,len,1);ntt(B,len,1);
    for(int i=0;i<len;++i)</pre>
      B[i]=(11)B[i]*(2-(11)t[i]*B[i]*mod+mod)*mod;
    ntt(B, len, -1); cls(B, l, len);
  void pmod(int*A){
    static int t[N];
    int l=k+1,len=1;while(len<=(k<<1))len<<=1;</pre>
    cpy(t, A, (k<<1)+1);
    reverse(t, t+(k << 1)+1);
    cls(t,1,len)
    ntt(t,len,1);
    for(int i=0;i<len;++i)t[i]=(11)t[i]*ib[i]%mod;</pre>
    ntt(t,len,-1);
    cls(t,1,len);
    reverse(t,t+1);
    ntt(t,len,1);
    for(int i=0;i<len;++i)t[i]=(11)t[i]*b[i]%mod;</pre>
    ntt(t,len,-1);
    cls(t,1,len);
    for(int i=0;i<k;++i)A[i]=(A[i]-t[i]+mod)%mod;</pre>
    cls(A,k,len);
  void pow(int*A,int n){
    if(n==1) {cls(A, 0, k+1); A[1]=1; return;}
    pow(A, n>>1);
    int len=1; while(len<=(k<<1))len<<=1;</pre>
    ntt(A,len,1);
    for(int i=0;i<len;++i)A[i]=(11)A[i]*A[i]%mod;</pre>
    ntt(A,len,-1);
    pmod(A);
    if(n&1){
       for(int i=k;i;--i)A[i]=A[i-1];A[0]=0;
      pmod(A);
  }
}
int main(){
  n=rd();k=rd();
  for(int i=1;i<=k;++i)f[i]=(mod+rd())%mod;</pre>
  for(int i=0;i<k;++i)h[i]=(mod+rd())%mod;</pre>
  for(int i=a[k]=b[k]=1;i<=k;++i)</pre>
    a[k-i]=b[k-i]=(mod-f[i])%mod;
  int len=1; while(len<=(k<<1))len<<=1;</pre>
  reverse(a,a+k+1);
  poly::inv(a,ib,len);
  poly::cls(ib,k+1,len);
  poly::ntt(b,len,1);
  poly::ntt(ib,len,1);
  poly::pow(a,n);
  int ans=0;
  for(int i=0;i<k;++i)inc(ans,(ll)a[i]*h[i]%mod);</pre>
  printf("%d\n",ans);
  return 0;
```

#### 5.13 Chinese Remainder

0 <= x < lcm(m1, m2)

```
lld crt(lld ans[], lld pri[], int n){
  lld M = 1, ret = 0;
  for(int i=0;i<n;i++) M *= pri[i];</pre>
  for(int i=0;i<n;i++){</pre>
    lld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
    ret += (ans[i]*(M/pri[i])%M * iv)%M;
    ret %= M;
  return ret;
}
/*
Another:
x = a1 \% m1
x = a2 \% m2
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
```

## 5.14 Berlekamp Massey

| \*/

```
// x: 1-base, p[]: 0-base
template<size_t N>
vector<llf> BM(llf x[N], size_t n){
  size_t f[N] = \{0\}, t = 0; llf d[N];
  vector<llf> p[N];
  for(size_t i=1,b=0;i<=n;++i) {</pre>
    for(size_t j=0;j<p[t].size();++j)</pre>
      d[i]+=x[i-j-1]*p[t][j];
    if(abs(d[i]-=x[i])<=EPS)continue;</pre>
    f[t]=i;if(!t){p[++t].resize(i);continue;}
    vector<llf> cur(i-f[b]-1);
    llf k=-d[i]/d[f[b]];cur.PB(-k);
    for(size_t j=0;j<p[b].size();j++)</pre>
      cur.PB(p[b][j]*k);
    if(cur.size()<p[t].size())cur.resize(p[t].size());</pre>
    for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];</pre>
    if(i-f[b]+p[b].size()>=p[t].size()) b=t;
    p[++t]=cur;
  return p[t];
```

```
5.15 NTT
// Remember coefficient are mod P
/* p=a*2^n+1
        2^n
                                        root
        65536
                      65537
   16
                                        3 */
        1048576
                      7340033
                                  7
   20
// (must be 2<sup>k</sup>)
template<LL P, LL root, int MAXN>
struct NTT{
  static LL bigmod(LL a, LL b) {
    LL res = 1;
    for (LL bs = a; b; b >>= 1, bs = (bs * bs) % P)
      if(b&1) res=(res*bs)%P;
    return res;
  static LL inv(LL a, LL b) {
    if(a==1)return 1;
    return (((LL)(a-inv(b%a,a))*b+1)/a)%b;
 LL omega[MAXN+1];
 NTT() {
    omega[0] = 1;
    LL r = bigmod(root, (P-1)/MAXN);
    for (int i=1; i<=MAXN; i++)</pre>
      omega[i] = (omega[i-1]*r)%P;
  // n must be 2^k
  void tran(int n, LL a[], bool inv_ntt=false){
    int basic = MAXN / n , theta = basic;
    for (int m = n; m >= 2; m >>= 1) {
      int mh = m >> 1;
      for (int i = 0; i < mh; i++) {</pre>
        LL w = omega[i*theta%MAXN];
        for (int j = i; j < n; j += m) {
  int k = j + mh;</pre>
          LL x = a[j] - a[k];
           if (x < 0) x += P;
           a[j] += a[k];
          if (a[j] > P) a[j] -= P;
          a[k] = (w * x) % P;
      theta = (theta * 2) % MAXN;
    for (int j = 1; j < n - 1; j++) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
    if (inv_ntt) {
      L\hat{L} ni = inv(n,P);
      reverse( a+1 , a+n );
      for (i = 0; i < n; i++)
        a[i] = (a[i] * ni) % P;
```

```
}
};
const LL P=2013265921, root=31;
const int MAXN=4194304;
NTT<P, root, MAXN> ntt;
```

## 5.16 Polynomial Operations

```
using VI = vector<int>;
Poly Inverse(Poly f) {
  int n = f.size()
  Poly q(1, fpow(f[0], kMod - 2));
  for (int s = 2;; s <<= 1) {
    if (f.size() < s) f.resize(s);</pre>
    Poly fv(f.begin(), f.begin() + s);
    Poly fq(q.begin(), q.end())
    fv.resize(s + s); fq.resize(s + s);
    ntt::Transform(fv, s + s);
    ntt::Transform(fq, s + s);
    for (int i = 0; i < s + s; ++i)
      fv[i] = 1LL * fv[i] * fq[i]%kMod * fq[i]%kMod;
    ntt::InverseTransform(fv, s + s);
    Poly res(s):
    for (int i = 0; i < s; ++i) {</pre>
      res[i] = kMod - fv[i];
      if (i < (s >> 1)) {
        int v = 2 * q[i] % kMod;
        (res[i] += v) >= kMod ? res[i] -= kMod : 0;
      }
    q = res;
    if (s >= n) break;
  q.resize(n);
  return q;
Poly Divide(const Poly &a, const Poly &b) {
  int n = a.size(), m = b.size(), k = 2;
  while (k < n - m + 1) k <<= 1;
  Poly ra(k), rb(k);
  for (int i = 0; i < min(n, k); ++i) ra[i] = a[n-1-i];
  for (int i = 0; i < min(m, k); ++i) rb[i] = b[m-1-i];
  auto rbi = Inverse(rb);
  auto res = Multiply(rbi, ra);
  res.resize(n - m + 1);
  reverse(res.begin(), res.end());
  return res:
Poly Modulo(const Poly &a, const Poly &b) {
  if (a.size() < b.size()) return a;</pre>
  auto dv = Multiply(Divide(a, b), b);
  assert(dv.size() == a.size());
  for (int i = 0; i < dv.size(); ++i)</pre>
    dv[i] = (a[i] + kMod - dv[i]) % kMod;
  while (!dv.empty() && dv.back() == 0) dv.pop_back();
  return dv;
Poly Integral(const Poly &f) {
  int n = f.size();
  VI res(n + 1);
  for (int i = 0; i < n; ++i)
    res[i+1] = 1LL * f[i] * fpow(i + 1, kMod - 2)%kMod;
  return res;
Poly Evaluate(const Poly &f, const VI &x) {
  if (x.empty()) return Poly();
  int n = x.size();
  vector<Poly> up(n * 2);
  for (int i = 0; i < n; ++i) up[i+n] = {kMod-x[i], 1};
  for (int i = n - 1; i > 0; --i)
    up[i] = Multiply(up[i * 2], up[i * 2 + 1]);
  vector<Poly> down(n * 2)
  down[1] = Modulo(f, up[1]);
  for (int i = 2; i < n * 2; ++i)
    down[i] = Modulo(down[i >> 1], up[i]);
  VI y(n);
  for (int i = 0; i < n; ++i) y[i] = down[i + n][0];</pre>
  return y;
Poly Interpolate(const VI &x, const VI &y) {
```

```
int n = x.size():
  vector<Poly> up(n * 2);
 for (int i = 0; i < n; ++i) up[i+n] = {kMod-x[i], 1};
for (int i = n - 1; i > 0; --i)
    up[i] = Multiply(up[i * 2], up[i * 2 + 1]);
  VI a = Evaluate(Derivative(up[1]), x);
  for (int i = 0; i < n; ++i)
    a[i] = 1LL * y[i] * fpow(a[i], kMod - 2) % kMod;
  vector<Poly> down(n * 2);
  for (int i = 0; i < n; ++i) down[i + n] = {a[i]};
  for (int i = n - 1; i > 0; --i) {
  auto lhs = Multiply(down[i * 2], up[i * 2 + 1]);
    auto rhs = Multiply(down[i * 2 + 1], up[i * 2]);
    assert(lhs.size() == rhs.size());
    down[i].resize(lhs.size());
    for (int j = 0; j < lhs.size(); ++j)</pre>
      down[i][j] = (lhs[j] + rhs[j]) % kMod;
  return down[1];
Poly Log(Poly f)
 int n = f.size();
  if (n == 1) return {0};
  auto d = Derivative(f);
 f.resize(n - 1);
  d = Multiply(d, Inverse(f));
  d.resize(n - 1);
  return Integral(d);
Poly Exp(Poly f) {
  int n = f.size();
  Poly q(1, 1); f[0] += 1;
  for (int s = 1; s < n; s <<= 1) {
  if (f.size() < s + s) f.resize(s + s);</pre>
    Poly g(f.begin(), f.begin() + s + s);
    Poly h(q.begin(), q.end())
    h.resize(s + s); h = Log(h);
    for (int i = 0; i < s + s; ++i)
      g[i] = (g[i] + kMod - h[i]) \% kMod;
    g = Multiply(g, q);
    g.resize(s + s); q = g;
  assert(q.size() >= n);
  q.resize(n);
  return q;
Poly SquareRootImpl(Poly f) {
  if (f.empty()) return {0};
  int z = QuadraticResidue(f[0], kMod), n = f.size();
  constexpr int kInv2 = (kMod + 1) >> 1;
  if (z == -1) return {-1};
  VI q(1, z);
  for (int s = 1; s < n; s <<= 1) {
    if (f.size() < s + s) f.resize(s + s);
    VI fq(q.begin(), q.end());
    fq.resize(s + s);
    VI f2 = Multiply(fq, fq);
    f2.resize(s + s);
    for (int i = 0; i < s + s; ++i)
      f2[i] = (f2[i] + kMod - f[i]) % kMod;
    f2 = Multiply(f2, Inverse(fq));
    f2.resize(s + s)
    for (int i = 0; i < s + s; ++i)
      fq[i] = (fq[i]+kMod - 1LL*f2[i]*kInv2%kMod)%kMod;
    q = fq;
  q.resize(n);
  return q;
Poly SquareRoot(Poly f) {
 int n = f.size(), m = 0;
  while (m < n \&\& f[m] == 0) m++;
  if (m == n) return VI(n);
  if (m & 1) return {-1}
  auto s = SquareRootImpl(VI(f.begin() + m, f.end()));
  if (s[0] == -1) return {-1};
 VI res(n);
  for (int i = 0; i < s.size(); ++i) res[i + m/2]=s[i];</pre>
  return res;
```

#### 5.17 FWT

```
/* xor convolution:
  * x = (x0, x1) , y = (y0, y1)
  *z = (x0y0 + x1y1 , x0y1 + x1y0 )
  * =>
  * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
  *z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
  *z = (1/2) *z'
  * or convolution:
  * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
  * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */const LL MOD = 1e9+7;
 inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
   for( int d = 1 ; d < N ; d <<= 1 ) {
     int d2 = d << 1
     for( int s = 0 ; s < N ; s += d2 )
       for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
  LL ta = x[ i ] , tb = x[ j ];</pre>
         x[i] = ta+tb;
         x[ j ] = ta-tb;
if( x[ i ] >= MOD ) x[ i ] -= MOD;
         if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
       }
   if( inv )
     for( int i = 0 ; i < N ; i++ ) {</pre>
       x[i] *= inv(N, MOD);
       x[ i ] %= MOD;
```

## 5.18 DiscreteLog

```
/ Baby-step Giant-step Algorithm
11d BSGS(11d P, 11d B, 11d N) {
  // find B^L = N mod P
  unordered_map<lld, int> R;
  1ld sq = (lld)sqrt(P);
  11d t = 1;
  for (int i = 0; i < sq; i++) {
    if (t == N) return i;
    if (!R.count(t)) R[t] = i;
    t = (t * B) % P;
  11d f = inverse(t, P);
  for(int i=0;i<=sq+1;i++) {</pre>
    if (R.count(N))
      return i * sq + R[N];
    N = (N * f) % P;
  return -1;
}
```

#### 5.19 Quadratic residue

if(p==2) return 1;

```
struct Status{
 11 x,y;
11 w:
Status mult(const Status& a,const Status& b,ll mod){
  res.x=(a.x*b.x+a.y*b.y%mod*w)%mod;
  res.y=(a.x*b.y+a.y*b.x)%mod;
  return res;
inline Status qpow(Status _base, 11 _pow, 11 _mod) {
  Status res = \{1, 0\};
  while(_pow>0){
    if(_pow&1) res=mult(res,_base,_mod);
    _base=mult(_base,_base,_mod);
    _pow>>=1;
  }
  return res;
inline 11 check(11 x,11 p){
 return qpow_mod(x,(p-1)>>1,p);
inline 11 get_root(11 n,11 p){
```

```
if(check(n,p)==p-1) return -1;
ll a;
while(true){
    a=rand()%p;
    w=((a*a-n)%p+p)%p;
    if(check(w,p)==p-1) break;
}
Status res = {a, 1}
res=qpow(res,(p+1)>>1,p);
return res.x;
}
5.20 De-Bruijn
```

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
 if (t > n) {
    if (n % p == 0)
     for (int i = 1; i <= p; ++i)
       res[sz++] = aux[i];
 } else {
    aux[t] = aux[t - p];
    db(t + 1, p, n, k);
    for (int i = aux[t - p] + 1; i < k; ++i) {
     aux[t] = i;
      db(t + 1, t, n, k);
 }
int de_bruijn(int k, int n) {
 // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
 if (k == 1) {
    res[0] = 0;
    return 1:
  for (int i = 0; i < k * n; i++) aux[i] = 0;
 sz = 0;
 db(1, 1, n, k);
 return sz;
```

#### 5.21 Simplex Construction

Standard form: maximize  $\sum_{1\leq i\leq n}c_ix_i$  such that for all  $1\leq j\leq m$ ,  $\sum_{1\leq i\leq n}A_{ji}x_i\leq b_j$  and  $x_i\geq 0$  for all  $1\leq i\leq n$ .

- 1. In case of minimization, let  $c_i^\prime = -c_i$
- 2.  $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- 3.  $\sum_{1 \le i \le n} A_{ji} x_i = b_j$ 
  - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$
  - $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

## 5.22 Simplex

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return vector<double>(n, -inf) if the solution doesn
     't exist
// return vector<double>(n, +inf) if the solution is
    unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9:
int n, m;
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
  double inv = 1.0 / d[r][s];
  for (int i = 0; i < m + 2; ++i) {
    for (int j = 0; j < n + 2; ++j) {
  if (i != r && j != s)
        d[i][j] = d[r][j] * d[i][s] * inv;
  }
```

```
for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
  d[r][s] = inv;
  swap(p[r], q[s]);
bool phase(int z) {
  int x = m + z;
  while (true) {
     int s = -1;
     for (int i = 0; i <= n; ++i) {
       if (!z && q[i] == -1) continue;
       if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
     if (d[x][s] > -eps) return true;
     int r = -1;
     for (int i = 0; i < m; ++i) {
       if (d[i][s] < eps) continue;
if (r == -1 || \</pre>
         d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
     if (r == -1) return false;
     pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
  m = b.size(), n = c.size();
  d = VVD(m + 2, VD(n + 2));
  for (int i = 0; i < m; ++i) {
     for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
  p.resize(m), q.resize(n + 1);
  for (int i = 0; i < m; ++i)
  p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];</pre>
  q[n] = -1, d[m + 1][n] = 1;
  int r = 0;
  for (int i = 1; i < m; ++i)
     if (d[i][n + 1] < d[r][n + 1]) r = i;
  if (d[r][n + 1] < -eps) {
     pivot(r, n);
     if (!phase(1) || d[m + 1][n + 1] < -eps)
     return VD(n, -inf);
for (int i = 0; i < m; ++i) if (p[i] == -1) {
       int s = min_element(d[i].begin(), d[i].end() - 1)
                 - d[i].begin();
       pivot(i, s);
  if (!phase(0)) return VD(n, inf);
  VD x(n);
  for (int i = 0; i < m; ++i)
     if (p[i] < n) x[p[i]] = d[i][n + 1];
   return x;
}}
```

# 6 Geometry

#### 6.1 Point Class

```
template<typename T>
struct Point{
  typedef long double llf:
  static constexpr llf EPS = 1e-8;
 T x, y
  Point(T _=0, T __=0): x(_), y(__){}
  template<typename T2>
    Point(const Point<T2>& a): x(a.x), y(a.y){}
  inline llf theta() const {
    return atan2((11f)y, (11f)x);}
  inline llf dis() const {
    return hypot((llf)x, (llf)y);}
  inline llf dis(const Point& o) const {
    return hypot((llf)(x-o.x), (llf)(y-o.y));}
  Point operator-(const Point& o) const {
    return Point(x-o.x, y-o.y);}
 Point operator+(const Point& o) const {
 return Point(x+o.x, y+o.y);}
Point operator*(const T& k) const {
    return Point(x*k, y*k);}
  Point operator/(const T& k) const {
    return Point(x/k, y/k);}
```

```
Point operator-() const {return Point(-x, -y);}
Point rot90() const {return Point(-y, x);}
template<typename T2>
bool in(const Circle<T2>& a) const {
  /* Add struct Circle at top */
  return a.o.dis(*this)+EPS <= a.r; }</pre>
bool equal(const Point& o, true_type) const {
  return fabs(x-o.x) < EPS and fabs(y-o.y) < EPS; }
bool equal(const Point& o, false_type) const {
return tie(x, y) == tie(o.x, o.y); }
bool operator==(const Point& o) const
  return equal(o, is_floating_point<T>()); }
bool operator!=(const Point& o) const {
  return !(*this == o); }
bool operator<(const Point& o) const {</pre>
  return theta() < o.theta();</pre>
  // sort like what pairs did
  // if (is_floating_point<T>())
       return fabs(x-o.x)<EPS?y<o.y:x<o.x;
  // else return tie(x, y) < tie(o.x, o.y);</pre>
friend inline T cross(const Point&a,const Point&b){
  return a.x*b.y - b.x*a.y; }
friend inline T dot(const Point& a, const Point &b){
  return a.x*b.x + a.y*b.y; }
friend ostream&operator<<(ostream&ss,const Point&o){</pre>
  ss<<"("<<o.x<<", "<<o.y<<")"; return ss; }
```

#### 6.2 Circle Class

```
template<typename T>
struct Circle{
    static constexpr llf EPS = 1e-8;
    Point<T> o; T r;
    vector<Point<llf>> operator&(const Circle& aa)const{
        // https://www.cnblogs.com/wangzming/p/8338142.html
        llf d=o.dis(aa.o);
        if(d > r+aa.r+EPS or d < fabs(r-aa.r)-EPS) return
            {};
        llf dt = (r*r - aa.r*aa.r)/d, d1 = (d+dt)/2;
        Point<llf>> dir = (aa.o-o); dir /= d;
        Point<llf>> pcrs = dir*d1 + o;
        dt=sqrt(max(0.0L, r*r - d1*d1)), dir=dir.rot90();
        return {pcrs + dir*dt, pcrs - dir*dt};
    }
};
```

## 6.3 Segment Class

```
const long double EPS = 1e-8;
template<typename T>
struct Segment{
  // p1.x < p2.x
 Line<T> base;
 Point<T> p1, p2;
 Segment(): base(Line<T>()), p1(Point<T>()), p2(Point<</pre>
    assert(on_line(p1, base) and on_line(p2, base));
 Segment(Line<T> _, Point<T> __, Point<T> __): base(_
), p1(__), p2(___){
    assert(on_line(p1, base) and on_line(p2, base));
  template<typename T2>
    Segment(const Segment<T2>& _): base(_.base), p1(_.
        p1), p2(_.p2) {}
  typedef Point<long double> Pt;
  friend bool on_segment(const Point<T>& p, const
      Segment& 1){
    if(on_line(p, 1.base))
      return (1.p1.x-p.x)*(p.x-1.p2.x)>=0 and (1.p1.y-p)
          .y)*(p.y-1.p2.y)>=0;
    return false;
  friend bool have_inter(const Segment& a, const
      Segment& b) {
    if(is_parallel(a.base, b.base)){
      return on_segment(a.p1, b) or on_segment(a.p2, b)
           or on_segment(b.p1, a) or on_segment(b.p2, a
          );
```

```
Pt inter = get_inter(a.base, b.base);
    return on_segment(inter, a) and on_segment(inter, b
  friend inline Pt get_inter(const Segment& a, const
      Segment& b){
    if(!have_inter(a, b)){
      return NOT_EXIST;
    }else if(is_parallel(a.base, b.base)){
      if(a.p1 == b.p1){
        if(on_segment(a.p2, b) or on_segment(b.p2, a))
            return INF_P;
        else return a.p1;
      else if(a.p1 == b.p2){
        if(on_segment(a.p2, b) or on_segment(b.p1, a))
            return INF_P
        else return a.p1;
      else if(a.p2 == b.p1){
        if(on_segment(a.p1, b) or on_segment(b.p2, a))
            return INF_P;
        else return a.p2;
      else if(a.p2 == b.p2){
        if(on_segment(a.p1, b) or on_segment(b.p1, a))
            return INF_P;
        else return a.p2;
      return INF_P;
    return get_inter(a.base, b.base);
  friend ostream& operator<<(ostream& ss, const Segment
      & o){
    ss<<o.base<<", "<<o.p1<<" ~ "<<o.p2;
    return ss;
template<typename T>
inline Segment<T> get_segment(const Point<T>& a, const
    Point<T>& b){
  return Segment<T>(get_line(a, b), a, b);
```

#### 6.4 Line Class

```
const Point<long double> INF_P(-1e20, 1e20):
const Point<long double> NOT_EXIST(1e20, 1e-20);
template<typename T>
struct Line{
  static constexpr long double EPS = 1e-8;
  // ax+by+c = 0
 T a, b, c;
Line(T _=0, T __=1, T ___=0): a(_), b(__), c(___){
assert(fabs(a)>EPS or fabs(b)>EPS);}
  template<typename T2>
    Line(const Line<T2>\&x): a(x.a), b(x.b), c(x.c){}
  typedef Point<long double> Pt;
  bool equal(const Line& o, true_type) const {
    return fabs(a-o.a)<EPS &&
    fabs(b-o.b)<EPS && fabs(c-o.b)<EPS;}</pre>
  bool equal(const Line& o, false_type) const {
    return a==o.a and b==o.b and c==o.c;}
  bool operator==(const Line& o) const {
    return equal(o, is_floating_point<T>());}
  bool operator!=(const Line& o) const {
    return !(*this == o);}
  friend inline bool on_line__(const Point<T>& p, const
       Line& 1, true_type){
    return fabs(1.a*p.x + 1.b*p.y + 1.c) < EPS;
  friend inline bool on_line__(const Point<T>& p, const
       Line& 1, false_type){
    return 1.a*p.x + 1.b*p.y + 1.c == 0;
  friend inline bool on_line(const Point<T>&p, const
      Line\& 1){
    return on_line__(p, l, is_floating_point<T>());
  friend inline bool is_parallel__(const Line& x, const
       Line& y, true_type){
    return fabs(x.a*y.b - x.b*y.a) < EPS;</pre>
```

```
friend inline bool is_parallel__(const Line& x, const
       Line& y, false_type){
    return x.a*y.b == x.b*y.a;
  friend inline bool is_parallel(const Line& x, const
      Line& y){
    return is_parallel__(x, y, is_floating_point<T>());
  friend inline Pt get_inter(const Line& x, const Line&
    typedef long double llf;
    if(x==y) return INF_P;
    if(is_parallel(x, y)) return NOT_EXIST;
    llf delta = x.a*y.b - x.b*y.a;
    llf delta_x = x.b*y.c - x.c*y.b;
    llf delta_y = x.c*y.a - x.a*y.c;
    return Pt(delta_x / delta, delta_y / delta);
  friend ostream& operator<<(ostream& ss, const Line& o
    ss<<o.a<<"x+"<<o.b<<"y+"<<o.c<<"=0";
    return ss;
template<typename T>
inline Line<T> get_line(const Point<T>& a, const Point<</pre>
    T > & b) {
  return Line<T>(a.y-b.y, b.x-a.x, (b.y-a.y)*a.x-(b.x-a
      .x)*a.y);
```

## 6.5 Triangle Circumcentre

#### 6.6 2D Convex Hull

```
template<typename T>
class ConvexHull_2D{
private:
  typedef Point<T> PT;
  vector<PT> d;
  struct myhash{
    uint64_t operator()(const PT& a) const {
      uint64_t xx=0, yy=0;
      memcpy(&xx, &a.x, sizeof(a.x));
      memcpy(&yy, &a.y, sizeof(a.y));
      uint64_t ret = xx*17+yy*31;
      ret = (\text{ret } ^{\prime} (\text{ret } >> 16))*0x9E3779B1;
      ret = (ret ^ (ret >> 13))*0xC2B2AE35;
      ret = ret ^ xx;
      return (ret ^ (ret << 3)) * yy;</pre>
  };
  unordered_set<PT, myhash> in_hull;
public:
  void init(){in_hull.clear();d.clear();}
  void insert(const PT& x){d.PB(x);}
  void solve(){
    sort(ALL(d), [](const PT& a, const PT& b){
      return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
    vector<PT> s(SZ(d)<<1); int o=0;
    for(auto p: d) {
      while(o \ge 2 \& cross(p-s[o-2], s[o-1]-s[o-2]) <= 0)
        0--;
```

```
s[o++] = p;
}
for(int i=SZ(d)-2, t = o+1;i>=0;i--){
    while(o>=t&&cross(d[i]-s[o-2],s[o-1]-s[o-2])<=0)
        o--;
    s[o++] = d[i];
}
s.resize(o-1); swap(s, d);
for(auto i: s) in_hull.insert(i);
}
vector<PT> get(){return d;}
bool in_it(const PT& x){
    return in_hull.find(x)!=in_hull.end();}
};
```

#### 6.7 2D Farthest Pair

```
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {
  while(abs(cross(stk[i+1]-stk[i],
        stk[(pos+1)%n]-stk[i])) >
        abs(cross(stk[i+1]-stk[i],
        stk[pos]-stk[i]))) pos = (pos+1)%n;
  ans = max({ans, dis(stk[i], stk[pos]),
        dis(stk[i+1], stk[pos]));
}
```

## 6.8 2D Closest Pair

```
struct Pt{
 11f x, y, d;
} arr[N];
inline llf dis(Pt a, Pt b){
  return hypot(a.x-b.x, a.y-b.y);
11f solve(){
  int cur = rand() % n;
  for(int i=0;i<n;i++) arr[i].d = dis(arr[cur], arr[i])</pre>
  sort(arr, arr+n, [](Pt a, Pt b){return a.d < b.d;});</pre>
  11f ans = 1e50;
  for(int i=0;i<n;i++){</pre>
    for(int j=i+1;j<n;j++){</pre>
      if(arr[j].d - arr[i].d > ans) break;
      ans = min(ans, dis(arr[i], arr[j]));
  }
  return ans;
```

## 6.9 Simulated Annealing

```
Ilf anneal() {
   mt19937 rnd_engine( seed );
   uniform_real_distribution< llf > rnd( 0, 1 );
   const llf dT = 0.001;
   // Argument p
   llf S_cur = calc( p ), S_best = S_cur;
   for ( llf T = 2000 ; T > EPS ; T -= dT ) {
        // Modify p to p_prime
        const llf S_prime = calc( p_prime );
        const llf delta_c = S_prime - S_cur;
        llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
        if ( rnd( rnd_engine ) <= prob )
            S_cur = S_prime, p = p_prime;
        if ( S_prime < S_best )
            S_best = S_prime, p_best = p_prime;
   }
   return S_best;
}</pre>
```

#### 6.10 Half Plane Intersection

```
inline int dcmp ( double x ) {
  if( fabs( x ) < eps ) return 0;
  return x > 0 ? 1 : -1;
}
```

```
struct Line {
  Point st, ed;
  double ang;
  Line(Point _s=Point(), Point _e=Point()):
   st(_s),ed(_e),ang(atan2(_e.y-_s.y,_e.x-_s.x)){}
  inline bool operator< ( const Line& rhs ) const {</pre>
    if(dcmp(ang - rhs.ang) != 0) return ang < rhs.ang;</pre>
    return dcmp( cross( st, ed, rhs.st ) ) < 0;</pre>
  }
};
// cross(pt, line.ed-line.st)>=0 <-> pt in half plane
vector< Line > lns;
deque< Line > que;
deque< Point > pt;
double HPI() {
  sort( lns.begin(), lns.end() );
  que.clear(); pt.clear();
que.push_back( lns[ 0 ] );
  for ( int i = 1 ; i < (int)lns.size() ; i ++ ) {</pre>
    if(!dcmp(lns[i].ang - lns[i-1].ang)) continue;
    while ( pt.size() > 0 &&
     dcmp(cross(lns[i].st,lns[i].ed,pt.back()))<0){</pre>
      pt.pop_back();que.pop_back();
    while ( pt.size() > 0 &&
     dcmp(cross(lns[i].st,lns[i].ed,pt.front()))<0){</pre>
      pt.pop_front(); que.pop_front();
    pt.push_back(get_point( que.back(), lns[ i ] ));
    que.push_back( lns[ i ] );
  while ( pt.size() > 0 &&
   dcmp(cross(que[0].st, que[0].ed, pt.back()))<0){</pre>
    que.pop_back();
    pt.pop_back();
  while ( pt.size() > 0 &&
   dcmp(cross(que.back().st,que.back().ed,pt[0]))<0){</pre>
    que.pop_front();
    pt.pop_front();
  pt.push_back(get_point(que.front(), que.back()));
  vector< Point > conv;
  for ( int i = 0 ; i < (int)pt.size() ; i ++ )</pre>
    conv.push_back( pt[ i ] );
  double ret = 0;
  for ( int i = 1 ; i + 1 < (int)conv.size() ; i ++ )</pre>
    ret += abs(cross(conv[0], conv[i], conv[i + 1]));
  return ret / 2.0;
}
```

## Ternary Search on Integer

```
int TernarySearch(int 1, int r) {
  // max value @ (1, r]
 while (r - 1 > 1){
    int m = (1 + r) >> 1;
    if (f(m) > f(m + 1)) r = m;
    else l = m;
 return 1+1;
```

## 6.12 Minimum Covering Circle

```
template<typename T>
Circle<llf> MinCircleCover(const vector<Point<T>>& pts)
  random_shuffle(ALL(pts));
  Circle<llf> c = \{pts[0], 0\};
  int n = SZ(pts);
  for(int i=0;i<n;i++){</pre>
    if(pts[i].in(c)) continue;
    c = \{pts[i], 0\};
    for(int j=0;j<i;j++){</pre>
      if(pts[j].in(c)) continue;
      c.o = (pts[i] + pts[j]) / 2;
      c.r = pts[i].dis(c.o);
      for(int k=0;k<j;k++){</pre>
        if(pts[k].in(c)) continue;
        c = get_circum(pts[i], pts[j], pts[k]);
```

```
}
}
return c:
```

#### KDTree (Nearest Point) 6.13

```
const int MXN = 100005;
struct KDTree
  struct Node {
    int x,y,x1,y1,x2,y2;
    int id,f;
Node *L, *R;
  } tree[MXN], *root;
  int n;
  LL dis2(int x1, int y1, int x2, int y2) {
    LL dx = x1-x2, dy = y1-y2;
    return dx*dx+dy*dy;
  static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
  static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
  void init(vector<pair<int,int>> ip) {
    n = ip.size();
    for (int i=0; i<n; i++) {
      tree[i].id = i;
      tree[i].x = ip[i].first;
      tree[i].y = ip[i].second;
    root = build_tree(0, n-1, 0);
  Node* build_tree(int L, int R, int d) {
    if (L>R) return nullptr
    int M = (L+R)/2; tree[M].f = d%2;
    nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
    tree[M].x1 = tree[M].x2 = tree[M].x;
    tree[M].y1 = tree[M].y2 = tree[M].y;
    tree[M].L = build_tree(L, M-1, d+1);
    if (tree[M].L) {
      tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
      tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
    tree[M].R = build_tree(M+1, R, d+1);
    if (tree[M].R) {
      tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
      tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
    return tree+M;
  int touch(Node* r, int x, int y, LL d2){
    LL dis = sqrt(d2)+1;
    if (x<r->x1-dis || x>r->x2+dis ||
        y<r->y1-dis || y>r->y2+dis)
      return 0;
    return 1;
  void nearest(Node* r,int x,int y,int &mID,LL &md2) {
    if (!r || !touch(r, x, y, md2)) return;
    LL d2 = dis2(r->x, r->y, x, y);
    if (d2 < md2 | | (d2 == md2 && mID < r->id)) {
      mID = r -> id;
      md2 = d2;
    // search order depends on split dim
    if ((r->f == 0 && x < r->x) ||
        (r->f == 1 && y < r->y))
      nearest(r->L, x, y, mID, md2);
      nearest(r->R, x, y, mID, md2);
    } else {
      nearest(r->R, x, y, mID, md2);
      nearest(r->L, x, y, mID, md2);
  int query(int x, int y) {
```

int id = 1029384756;

LL d2 = 102938475612345678LL;

nearest(root, x, y, id, d2);

```
7 Stringology
```

return id:

#### 7.1 Hash

|} tree;

```
class Hash{
private:
  const int p = 127, q = 1051762951;
int sz, prefix[N], power[N];
  int add(int x, int y){return x+y>=q?x+y-q:x+y;}
  int sub(int x, int y){return x-y<0?x-y+q:x-y;}</pre>
  int mul(int x, int y){return 1LL*x*y%q;}
public:
  void init(const string &x){
    sz = x.size();prefix[0]=0;power[0]=1;
    for(int i=1;i<=sz;i++)</pre>
      prefix[i]=add(mul(prefix[i-1], p), x[i-1]);
    for(int i=1;i<=sz;i++)power[i]=mul(power[i-1], p);</pre>
  int query(int 1, int r){
    // 1-base (1, r]
    return sub(prefix[r], mul(prefix[1], power[r-1]));
};
```

## 7.2 Suffix Array

```
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
int x[maxn], p[maxn], q[maxn * 2];
// sa[i]: sa[i]-th suffix is the \
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *sa, int *c, int n, int z) {
  memset(sa, 0, sizeof(int) * n);
  memcpy(x, c, sizeof(int) * z);
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
 memcpy(x + 1, c, sizeof(int) * (z - 1));
  for (int i = 0; i < n; ++i)
    if (sa[i] && !t[sa[i] - 1])
sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  memcpy(x, c, sizeof(int) * z);
  for (int i = n - 1; i >= 0; --i)
  if (sa[i] && t[sa[i] - 1])
      sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q,
bool *t, int *c, int n, int z)
  bool uniq = t[n - 1] = true;
  int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
  memset(c, 0, sizeof(int) * z);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
  for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
    return:
  for (int i = n - 2; i >= 0; --i)
    t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i)
    if (t[i] && !t[i - 1])
      sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
  for (int i = 0; i < n; ++i) {
    if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
    bool neq = last < 0 || \</pre>
     memcmp(s + sa[i], s + last,
  (p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
    ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
  pre(sa, c, n, z);
```

```
for (int i = nn - 1:
                           i >= 0: --i)
     sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
void build(const string &s) {
  for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
  _s[(int)s.size()] = 0; // s shouldn't contain 0 sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
  for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];</pre>
  for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;
int ind = 0; hi[0] = 0;</pre>
  for (int i = 0; i < (int)s.size(); ++i) {</pre>
     if (!rev[i]) {
       ind = 0;
       continue;
     while (i + ind < (int)s.size() && \</pre>
      s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
     hi[rev[i]] = ind ? ind-- : 0;
}}
```

## 7.3 Aho-Corasick Algorithm

```
class AhoCorasick{
  private:
    static constexpr int Z = 26;
    struct node{
      node *nxt[ Z ], *fail;
      vector< int > data;
      node(): fail( nullptr ) {
        memset( nxt, 0, sizeof( nxt ) );
        data.clear();
    } *rt;
    inline int Idx( char c ) { return c - 'a'; }
  public:
    void init() { rt = new node(); }
    void add( const string& s, int d ) {
      node* cur = rt;
      for ( auto c : s ) {
        if ( not cur->nxt[ Idx( c ) ] )
          cur->nxt[ Idx( c ) ] = new node();
        cur = cur->nxt[ Idx( c ) ];
      cur->data.push_back( d );
    void compile() {
  vector< node* > bfs;
      size_t ptr = 0;
      for ( int i = 0 ; i < Z ; ++ i ) {
        if ( not rt->nxt[ i ] )
          continue;
        rt->nxt[ i ]->fail = rt;
        bfs.push_back( rt->nxt[ i ] );
      while ( ptr < bfs.size() ) {</pre>
        node* u = bfs[ ptr ++ ];
for ( int i = 0 ; i < Z ; ++ i ) {</pre>
          if ( not u->nxt[ i ] )
            continue;
          node* u_f = u->fail;
          while ( u_f ) {
            if ( not u_f->nxt[ i ] ) {
               u_f = u_f->fail; continue;
            u->nxt[ i ]->fail = u_f->nxt[ i ];
            break;
          if ( not u_f ) u->nxt[ i ]->fail = rt;
          bfs.push_back( u->nxt[ i ] );
        }
      }
    void match( const string& s, vector< int >& ret ) {
      node* u = rt;
      for ( auto c : s ) {
        while ( u != rt and not u->nxt[ Idx( c ) ] )
          u = u->fail;
        u = u - nxt[Idx(c)];
        if ( not u ) u = rt;
```

```
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        node* tmp = u;
        while ( tmp != rt ) {
          for ( auto d : tmp->data )
           ret.push_back( d );
                                                         }
          tmp = tmp->fail;
      }
    }
} ac;
      Suffix Automaton
7.4
struct Node{
  Node *green, *edge[26];
                                                         }
  int max_len;
  Node(const int _max_len)
    : green(NULL), max_len(_max_len){
    memset(edge,0,sizeof(edge));
} *ROOT, *LAST;
void Extend(const int c) {
  Node *cursor = LAST;
```

```
LAST = new Node((LAST->max_len) + 1);
 for(;cursor&&!cursor->edge[c]; cursor=cursor->green)
   cursor->edge[c] = LAST;
 if (!cursor)
   LAST->green = ROOT;
 else {
   Node *potential_green = cursor->edge[c];
    if((potential_green->max_len)==(cursor->max_len+1))
     LAST->green = potential_green;
    else {
//assert(potential_green->max_len>(cursor->max_len+1));
     Node *wish = new Node((cursor->max_len) + 1);
      for(;cursor && cursor->edge[c]==potential_green;
           cursor = cursor->green)
        cursor->edge[c] = wish;
      for (int i = 0; i < 26; i++)
        wish->edge[i] = potential_green->edge[i];
      wish->green = potential_green->green;
     potential_green->green = wish;
      LAST->green = wish;
 }
char S[10000001], A[10000001];
int N;
int main(){
  scanf("%d%s", &N, S);
  ROOT = LAST = new Node(0);
  for (int i = 0; S[i]; i++)
    Extend(S[i] - 'a');
 while (N--){
    scanf("%s", A);
    Node *cursor = ROOT;
    bool ans = true;
    for (int i = 0; A[i]; i++){
     cursor = cursor->edge[A[i] - 'a'];
     if (!cursor) {
        ans = false;
        break;
     }
   puts(ans ? "Yes" : "No");
  return 0:
```

## 7.5 KMP

```
vector<int> kmp(const string &s) {
 vector<int> f(s.size(), 0);
  /* f[i] = length of the longest prefix
     (excluding s[0:i]) such that it coincides
     with the suffix of s[0:i] of the same length */
  /* i + 1 - f[i] is the length of the
     smallest recurring period of s[0:i] */
  int k = 0:
  for (int i = 1; i < (int)s.size(); ++i) {</pre>
   while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
    if (s[i] == s[k]) ++k;
```

```
f[i] = k;
  return f;
vector<int> search(const string &s, const string &t) {
  // return 0-indexed occurrence of t in s
  vector<int> f = kmp(t), r;
for (int i = 0, k = 0; i < (int)s.size(); ++i)</pre>
    while(k > 0 && (k==(int)t.size() || s[i]!=t[k]))
      k = f[k - 1];
    if (s[i] == t[k]) ++k;
    if (k == (int)t.size()) r.push_back(i-t.size()+1);
  return res;
```

#### 7.6 Z value

```
char s[MAXN];
int len,z[MAXN];
void Z_value() {
  int i,j,left,right;
  left=right=0; z[0]=len;
  for(i=1;i<len;i++) {</pre>
    j=max(min(z[i-left], right-i),0);
    for(;i+j<len&&s[i+j]==s[j];j++);
    z[i]=j;
    if(i+z[i]>right) {
      right=i+z[i];
      left=i;
  }
}
```

#### 7.7 Manacher

```
int z[maxn];
int manacher(const string& s) {
  string t = ".";
  for(char c:s)) t += c, t += '.';
  int 1 = 0, r = 0, ans = 0;
  for (int i = 1; i < t.length(); ++i) {</pre>
    z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
while (i - z[i] >= 0 && i + z[i] < t.length()) {
       if(t[i - z[i]] == t[i + z[i]]) ++z[i];
       else break;
    if (i + z[i] > r) r = i + z[i], l = i;
  for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
  return ans;
```

#### 7.8 Lexico Smallest Rotation

```
string mcp(string s){
  int n = s.length();
  s += s;
  int i=0, j=1;
  while (i<n && j<n){</pre>
    int k = 0:
    while (k < n \&\& s[i+k] == s[j+k]) k++;
    if (s[i+k] \le s[j+k]) j += k+1;
    else i += k+1;
    if (i == j) j++;
  int ans = i < n ? i : j;</pre>
  return s.substr(ans, n);
}
```

#### 7.9 BWT

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a'
  vector<int> v[ SIGMA ];
  void BWT(char* ori, char* res){
    // make ori -> ori + ori
    // then build suffix array
```

```
}
void iBWT(char* ori, char* res){
    for( int i = 0 ; i < SIGMA ; i ++ )
        v[ i ].clear();
    int len = strlen( ori );
    for( int i = 0 ; i < len ; i ++ )
        v[ ori[i] - BASE ].push_back( i );
    vectorsint> a;
    for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )
        for( auto j : v[ i ] ){
            a.push_back( j );
            ori[ ptr ++ ] = BASE + i;
        }
    for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
        res[ i ] = ori[ a[ ptr ] ];
        ptr = a[ ptr ];
    }
    res[ len ] = 0;
}
bwt;</pre>
```

### 7.10 Palindromic Tree

```
struct palindromic_tree{
  struct node{
    int next[26],f,len;
    int cnt, num, st, ed;
    node(int l=0):f(0),len(1),cnt(0),num(0)  {
      memset(next, 0, sizeof(next)); }
 vector<node> st;
 vector<char> s;
  int last,n;
 void init(){
    st.clear();s.clear();last=1; n=0;
    st.push_back(0);st.push_back(-1);
    st[0].f=1;s.push_back(-1); }
  int getFail(int x){
    while(s[n-st[x].len-1]!=s[n])x=st[x].f;
    return x;}
  void add(int c){
    s.push_back(c-='a'); ++n;
    int cur=getFail(last);
    if(!st[cur].next[c]){
      int now=st.size();
      st.push_back(st[cur].len+2);
      st[now].f=st[getFail(st[cur].f)].next[c];
st[cur].next[c]=now;
      st[now].num=st[st[now].f].num+1;
    last=st[cur].next[c];
    ++st[last].cnt;}
  int size(){ return st.size()-2;}
int main() {
  string s; cin >> s; pt.init();
  for (int i=0; i<SZ(s); i++) {
    int prvsz = pt.size(); pt.add(s[i]);
    if (prvsz != pt.size()) {
   int r = i, l = r - pt.st[pt.last].len + 1;
      // pal @ [1,r]: s.substr(1, r-l+1)
    }
  return 0:
```

## 8 Misc

#### 8.1 Theorems

#### 8.1.1 Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii}=d(i)$ ,  $L_{ij}=-c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|{\rm det}(\tilde{L}_{rr})|$  .

## 8.1.2 Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

### 8.1.3 Cayley's Formula

- Given a degree sequence  $d_1,d_2,\ldots,d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .

#### 8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all  $1 \le k \le n$ .

#### 8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

#### 8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let  $N_G(W)$  denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff  $\forall W\subseteq X, |W|\leq |N_G(W)|$ 

## 8.1.7 Euler's planar graph formula

V - E + F = C + 1,  $E \le 3V - 6$ (?)

#### 8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have  $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$ 

## 8.2 MaximumEmptyRect

```
int max_empty_rect(int n, int m, bool blocked[N][N]){
  static int mxu[2][N], me=0,he=1,ans=0;
  for(int i=0;i<m;i++) mxu[he][i]=0;</pre>
  for(int i=0;i<n;i++){</pre>
    stack<PII, vector<PII>> stk;
    for(int j=0;j<m;++j){</pre>
      if(blocked[i][j]) mxu[me][j]=0;
      else mxu[me][j]=mxu[he][j]+1;
      int la = j;
      while(!stk.empty()&&stk.top().FF>mxu[me][j]){
        int x1 = i - stk.top().FF, x2 = i;
        int y1 = stk.top().SS, y2 = j;
        la = stk.top().SS; stk.pop();
        ans=\max(ans, (x2-x1)*(y2-y1));
      if(stk.empty()||stk.top().FF<mxu[me][j])</pre>
        stk.push({mxu[me][j],la});
    while(!stk.empty()){
      int x1 = i - stk.top().FF, x2 = i;
      int y1 = stk.top().SS-1, y2 = m-1;
      stk.pop();
      ans=max(ans, (x2-x1)*(y2-y1));
    swap(me,he);
  }
  return ans;
}
```

### 8.3 DP-opt Condition

## 8.3.1 totally monotone (concave/convex)

 $\begin{array}{ll} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \Longrightarrow B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \Longrightarrow B[i][j'] \geq B[i'][j'] \end{array}$ 

### 8.3.2 monge condition (concave/convex)

 $\begin{array}{l} \forall i < i', j < j', \ B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', \ B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}$ 

## 8.4 Convex 1D/1D DP

```
struct segment {
             int i, l, r
              segment() {}
              segment(int a, int b, int c): i(a), l(b), r(c) {}
  inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
  void solve() {
             dp[0] = 0;
              deque<segment> dq; dq.push_back(segment(0, 1, n));
              for (int i = 1; i <= n; ++i) {
                        dp[i] = f(dq.front().i, i);
                        while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
                        dq.front().l = i + 1
                         segment seg = segment(i, i + 1, n);
                       while (dq.size() &&
                                    f(i, \dot{d}, \dot{d}, back(\dot{d}, 1) < f(\dot{d}, back(\dot{d}, \dot{d}, 
                                               dq.pop_back();
                        if (dq.size())
                                   int d = 1 \ll 20, c = dq.back().1;
                                    while (d \gg 1) if (c + d \ll d, back().r)
                                              if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
                                    dq.back().r = c; seg.l = c + 1;
                        if (seg.1 <= n) dq.push_back(seg);</pre>
}
```

## 8.5 ConvexHull Optimization

```
inline 1ld DivCeil(1ld n, 1ld d) { // ceil(n/d)
  return n / d + (((n < 0) != (d > 0)) && (n % d));
struct Line {
  static bool flag;
  11d a, b, 1, r; // y=ax+b in [1, r)
  11d operator()(11d x) const { return a * x + b; }
  bool operator<(const Line& i) const {</pre>
    return flag ? tie(a, b) < tie(i.a, i.b) : l < i.l;</pre>
  11d operator&(const Line& i) const {
    return DivCeil(b - i.b, i.a - a);
bool Line::flag = true;
class ConvexHullMax {
  set<Line> L;
public:
 ConvexHullMax() { Line::flag = true; }
  void InsertLine(lld a, lld b) { // add y = ax + b
    Line now = \{a, b, -INF, INF\};
    if (L.empty())
      L.insert(now);
      return;
    Line::flag = true;
    auto it = L.lower_bound(now);
    auto prv = it == L.begin() ? it : prev(it);
    if (it != L.end() && ((it != L.begin() &&
      (*it)(it->1) >= now(it->1) &&
      (*prv)(prv->r - 1) >= now(prv->r - 1)) ||
      (it == L.begin() && it->a == now.a))) return;
    if (it != L.begin())
      while (prv != L.begin() &&
        (*prv)(prv->1) \le now(prv->1))
          prv = --L.erase(prv);
      if (prv == L.begin() && now.a == prv->a)
        L.erase(prv);
    if (it != L.end())
      while (it != --L.end() &&
        (*it)(it->r) \le now(it->r))
          it = L.erase(it);
    if (it != L.begin()) {
      prv = prev(it);
      const_cast<Line*>(&*prv)->r=now.l=((*prv)&now);
    if (it != L.end())
      const_cast<Line*>(&*it)->l=now.r=((*it)&now);
    L.insert(it, now);
```

```
flld Query(lld a) const { // query max at x=a
   if (L.empty()) return -INF;
   Line::flag = false;
   auto it = --L.upper_bound({0, 0, a, 0});
   return (*it)(a);
};
```

## 8.6 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
   int s = 0;
   for (int i = 2; i <= n; i++)
        s = (s + m) % i;
   return s;
}
// died at kth
int kth(int n, int m, int k){
   if (m == 1) return n-1;
   for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
   return k;
}
```

## 8.7 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
  dfn[u]=low[u]=++dfs_idx;
  for(int i=0;i<(int)init_g[u].size();i++){</pre>
    int v=init_g[u][i];
    if(v==par[u]) continue;
    if(!dfn[v]){
      par[v]=u;
      tarjan(v);
      low[u]=min(low[u],low[v]);
      if(dfn[u]<low[v]){</pre>
        g[u].push_back(v)
        g[v].push_back(u);
    }else{
      low[u]=min(low[u],dfn[v]);
      if(dfn[v]<dfn[u]){</pre>
        int temp_v=u;
        bcc_id++;
        while(temp_v!=v){
          g[bcc_id+n].push_back(temp_v);
          g[temp_v].push_back(bcc_id+n);
          temp_v=par[temp_v];
        g[bcc_id+n].push_back(v);
        g[v].push_back(bcc_id+n);
        reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
    }
  }
int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
void dfs(int u,int fa){
  if(u<=n){
    for(int i=0;i<(int)g[u].size();i++){</pre>
      int v=g[u][i];
      if(v==fa) continue;
      dfs(v,u);
      memset(tp,0x8f,sizeof tp);
        tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
        tp[1]=max(
          dp[u][0]+dp[v][0]+1
          dp[u][1]+max(dp[v][0],dp[v][1])
      }else{
        tp[0]=dp[u][0]+dp[v][0];
        tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
      dp[u][0]=tp[0],dp[u][1]=tp[1];
  }else{
    for(int i=0;i<(int)g[u].size();i++){</pre>
```

for(int i = D[c]; i != c; i = D[i])

U[D[i]] = U[i];

D[U[j]] = D[j];

--S[col[j]];

for(int j = R[i]; j != i; j = R[j]) {

```
int v=g[u][i];
      if(v==fa) continue;
                                                                 void resume(int c) {
                                                                    L[R[c]] = c; R[L[c]] = c;
      dfs(v,u);
                                                                    for(int i = U[c]; i != c; i = U[i])
                                                                      for(int j = L[i]; j != i; j = L[j]) {
  U[D[j]] = j;
    min_dp[0][0]=0;
    min_dp[1][1]=1;
    \min_{dp[0][1]=\min_{dp[1][0]=-0}} x3f3f3f3f3f;
                                                                        D[U[j]] = j
    for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                        ++S[col[j]];
      int v=g[u][i];
                                                                    }
      if(v==fa) continue;
      memset(tmp,0x8f,sizeof tmp);
                                                                 void dance(int d) {
      tmp[0][0]=max(
                                                                    if(d>=ansd) return;
        \min_{dp[0][0]+\max(dp[v][0],dp[v][1]),
                                                                    if(R[0] == 0) {
        min_dp[0][1]+dp[v][0]
                                                                      ansd = d;
                                                                      return;
      tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
      tmp[1][0]=max(
                                                                    int c = R[0];
        min_dp[1][0]+max(dp[v][0],dp[v][1]),
                                                                    for(int i = R[0]; i; i = R[i])
        \min_{dp[1][1]+dp[v][0]}
                                                                      if(S[i] < S[c]) c = i;
                                                                    remove(c);
                                                                    for(int i = D[c]; i != c; i = D[i]) {
      tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
      memcpy(min_dp,tmp,sizeof tmp);
                                                                      ans[d] = row[i];
                                                                      for(int j = R[i]; j != i; j = R[j])
    dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
                                                                        remove(col[j]);
    dp[u][0]=min_dp[0][0];
                                                                      dance(d+1);
                                                                      for(int j = L[i]; j != i; j = L[j])
                                                                        resume(col[j]);
int main(){
  int m,a,b;
                                                                    resume(c);
  scanf("%d%d",&n,&m);
for(int i=0;i<m;i++){</pre>
                                                               } sol;
    scanf("%d%d",&a,&b);
    init_g[a].push_back(b);
                                                               8.9 Tree Knapsack
    init_g[b].push_back(a);
                                                               int dp[N][K];PII obj[N];
  par[1]=-1;
                                                               vector<int> G[N];
  tarjan(1);
                                                               void dfs(int u, int mx){
  dfs(1,-1);
                                                                 for(int s: G[u])
  printf("%d\n", max(dp[1][0], dp[1][1]));
                                                                    if(mx < obj[s].first) continue;</pre>
  return 0;
                                                                    for(int i=0;i<=mx-obj[s].FF;i++)</pre>
                                                                      dp[s][i] = dp[u][i];
                                                                    dfs(s, mx - obj[s].first);
                                                                    for(int i=obj[s].FF;i<=mx;i++)</pre>
8.8
      DLX
                                                                      dp[u][i] = max(dp[u][i],
                                                                        dp[s][i - obj[s].FF] + obj[s].SS);
struct DLX {
  const static int maxn=210;
  const static int maxm=210;
                                                               int main(){
  const static int maxnode=210*210;
                                                                 int n, k; cin >> n >> k;
  int n, m, size, row[maxnode], col[maxnode];
                                                                 for(int i=1;i<=n;i++){</pre>
  int U[maxnode], D[maxnode], L[maxnode], R[maxnode];
                                                                    int p; cin >> p;
  int H[maxn], S[maxm], ansd, ans[maxn];
                                                                    G[p].push_back(i);
  void init(int _n, int _m) {
                                                                    cin >> obj[i].FF >> obj[i].SS;
    n = _n, m = _m;
    for(int i = 0; i <= m; ++i) {</pre>
                                                                 dfs(0, k); int ans = 0;
      S[i] = 0
                                                                 for(int i=0;i<=k;i++) ans = max(ans, dp[0][i]);
      U[i] = D[i] = i;
                                                                 cout << ans << '\n';
      L[i] = i-1, R[i] = i+1;
                                                                 return 0;
    R[L[0] = size = m] = 0;
    for(int i = 1; i <= n; ++i) H[i] = -1;
                                                               8.10 N Queens Problem
  void Link(int r, int c) {
    ++S[col[++size] = c];
                                                               vector< int > solve( int n ) {
    row[size] = r; D[size] = D[c];
U[D[c]] = size; U[size] = c; D[c] = size;
                                                                 // no solution when n=2, 3
                                                                 vector< int > ret;
    if(H[r] < 0) H[r] = L[size] = R[size] = size;
                                                                 if ( n % 6 == 2 ) {
                                                                    for ( int i = 2 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
    else {
      R[size] = R[H[r]];
                                                                    ret.push_back( 3 ); ret.push_back( 1 );
for ( int i = 7 ; i <= n ; i += 2 )
  ret.push_back( i );
      L[R[H[r]]] = size;
      L[size] = H[r];
      R[H[r]] = size;
                                                                    ret.push_back( 5 );
                                                                 } else if ( n % 6 == 3 ) {
                                                                    for ( int i = 4 ; i <= n ; i += 2 )
  void remove(int c) {
    L[R[c]] = L[c]; R[L[c]] = R[c];
                                                                      ret.push_back( i );
```

ret.push\_back( 2 );

for ( int i = 5 ; i <= n ; i += 2 )
 ret.push\_back( i );</pre>

} else {
 for ( int i = 2 ; i <= n ; i += 2 )</pre>

ret.push\_back( 1 ); ret.push\_back( 3 );

```
ret.push_back( i );
for ( int i = 1 ; i <= n ; i += 2 )
    ret.push_back( i );
}
return ret;
}</pre>
```