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Basic
 vimrc
se is nu rnu bs=2 ru mouse=a encoding=utf-8
se cin et sw=4 sts=4 t_Co=256 tgc sc hls ls=2
svn on
colorscheme desert
filetype indent on
inoremap {<CR>} {<CR>} < ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 -
 DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
 Wconversion -fsanitize=address -fsanitize=undefined
  -g && echo success<CR>
```

map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 -DKISEKI && echo success<CR>

1.2 Increase Stack

map <F10> <ESC>:!./"%<"<CR>

```
const int size = 256 << 20;</pre>
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));
```

1.3 Pragma Optimization

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,sse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
```

1.4 IO Optimization

```
static inline int gc() {
  static char buf[ 1 << 20 ], *p = buf, *end = buf;</pre>
 if ( p == end ) {
  end = buf + fread( buf, 1, 1 << 20, stdin );
   if ( end == buf ) return EOF;
  p = buf;
 return *p++;
template < typename T >
static inline bool gn( T &_ ) {
  register int c = gc(); register T __ = 1; _ = 0;
  while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
 if(c == '-') { __ = -1; c = gc(); }
if(c == EOF) return false;
 while('0'<=c&&c<='9') _{-} = _{-} * 10 + c _{-} '0', c = gc();
  _ *= __;
 return true;
```

```
template < typename T, typename ...Args >
static inline bool gn( T &x, Args &...args )
{ return gn(x) && gn(args...); }
```

2 Data Structure

2.1 Dark Magic

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using __gnu_pbds::pairing_heap_tag;
using __gnu_pbds::binary_heap_tag;
using __gnu_pbds::binomial_heap_tag;
using __gnu_pbds::rc_binomial_heap_tag;
using __gnu_pbds::thin_heap_tag;
template<typename T>
using pbds_heap=__gnu_pbds::prioity_queue<T, less<T>, \
                   pairing_heap_tag>;
// __gnu_pbds::priority_queue<T,less<T>>::
    point_iterator
// x = pq.push(10); pq.modify(x, 87); a.join(b);
using __gnu_pbds::rb_tree_tag;
using __gnu_pbds::ov_tree_tag;
using __gnu_pbds::splay_tree_tag;
template<typename T>
using ordered_set = __gnu_pbds::tree<T,\</pre>
__gnu_pbds::null_type,less<T>,rb_tree_tag,\
 _gnu_pbds::tree_order_statistics_node_update>;
// find_by_order, order_of_key
template<typename A, typename B>
using hTable1=__gnu_pbds::cc_hash_table<A,B>;
template<typename A, typename B>
using hTable2=__gnu_pbds::gp_hash_table<A,B>;
```

2.2 Disjoint Set

```
class DJS {
private:
vector< int > fa, sz, sv;
vector< pair< int*, int > > opt;
void assign( int *k, int v ) {
 opt.emplace_back( k, *k );
  *k = v;
public:
void init( int n ) {
 fa.resize( n ); iota( fa.begin(), fa.end(), 0 );
sz.resize( n ); fill( sz.begin(), sz.end(), 1 );
 opt.clear();
 int query(int x) {return fa[x] == x?x:query(fa[x]);}
void merge( int a, int b ) -
 int af = query( a ), bf = query( b );
 if( af == bf ) return;
if( sz[ af ] < sz[ bf ] ) swap( af, bf );</pre>
 assign( &fa[ bf ], fa[ af ] );
 assign( &sz[ af ], sz[ af ] + sz[ bf ] );
void save() { sv.push_back( (int) opt.size() ); }
void undo() {
  int ls = sv.back(); sv.pop_back();
 while ( ( int ) opt.size() > ls )
   pair< int*, int > cur = opt.back();
   *cur.first = cur.second;
   opt.pop_back();
};
```

2.3 Link-Cut Tree

```
struct Node{
Node *par,*ch[2];
int xor_sum,v;
bool is_rev;
Node(int _v){
  v=xor_sum=_v;is_rev=false;
  par=ch[0]=ch[1]=nullptr;
}
inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
inline void down(){
  if(is_rev){
    if(ch[0]!=nullptr) ch[0]->set_rev();
    if(ch[1]!=nullptr) ch[1]->set_rev();
```

```
is_rev=false;
 inline void up(){
  xor_sum=v;
  if(ch[0]!=nullptr){
   xor_sum^=ch[0]->xor_sum;
   ch[0]->par=this;
  if(ch[1]!=nullptr){
   xor_sum^=ch[1]->xor_sum;
   ch[1]->par=this;
  }
 inline bool is_root(){
  return par==nullptr ||\
   (par->ch[0]!=this && par->ch[1]!=this);
 bool is_rch(){return !is_root() && par->ch[1]==this;}
} *node[maxn],*stk[maxn];
int top:
void to_child(Node* p,Node* c,bool dir){
p->ch[dir]=c;
p->up();
inline void rotate(Node* node){
 Node* par=node->par;
 Node* par_par=par->par
 bool dir=node->is_rch()
 bool par_dir=par->is_rch();
 to_child(par, node->ch[!dir], dir);
 to_child(node,par,!dir);
 if(par_par!=nullptr && par_par->ch[par_dir]==par)
 to_child(par_par, node, par_dir);
 else node->par=par_par;
inline void splay(Node* node){
Node* tmp=node;
 stk[top++]=node;
 while(!tmp->is_root()){
  tmp=tmp->par;
  stk[top++]=tmp;
 while(top) stk[--top]->down();
 for(Node *fa=node->par;
  !node->is_root();
  rotate(node),fa=node->par)
  if(!fa->is_root())
   rotate(fa->is_rch()==node->is_rch()?fa:node);
inline void access(Node* node){
 Node* last=nullptr;
 while(node!=nullptr){
  splay(node)
  to_child(node, last, true);
 last=node:
 node=node->par;
inline void change_root(Node* node){
access(node);splay(node);node->set_rev();
inline void link(Node* x, Node* y){
change_root(x);splay(x);x->par=y;
inline void split(Node* x, Node* y){
change_root(x);access(y);splay(x)
 to_child(x,nullptr,true);y->par=nullptr;
inline void change_val(Node* node,int v){
access(node);splay(node);node->v=v;node->up();
inline int query(Node* x,Node* y){
 change_root(x);access(y);splay(y);
 return y->xor_sum;
inline Node* find_root(Node* node){
 access(node);splay(node);
 Node* last=nullptr;
 while(node!=nullptr){
 node->down();last=node;node=node->ch[0];
```

```
return last;
                                                               } else {
                                                                R->lc = merge( L, R->lc ); R->pull();
set<pii> dic;
                                                                return R;
inline void add_edge(int u,int v){
if(u>v) swap(u,v)
 if(find_root(node[u])==find_root(node[v])) return;
                                                              void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                               if ( not rt ) L = R = nullptr;
dic.insert(pii(u,v));
link(node[u],node[v]);
                                                               else if( sz( rt->lc ) + 1 <= k ) {
                                                                L = rt
inline void del_edge(int u,int v){
                                                                split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
if(u>v) swap(u,v);
                                                                L->pull();
 if(dic.find(pii(u,v))==dic.end()) return;
                                                               } else {
dic.erase(pii(u,v))
                                                                R = rt:
split(node[u],node[v]);
                                                                split_by_size( rt->lc, k, L, R->lc );
                                                                R->pull();
2.4 LiChao Segment Tree
struct Line{
                                                              #undef sz
int m, k, id;
Line() : id( -1 ) {}
                                                             2.6 Sparse Table
Line( int a, int b, int c )
  : m( a ), k( b ), id( c ) {}
                                                             template < typename T, typename Cmp_ = less< T > >
 int at( int x ) { return m * x + k; }
                                                             class SparseTable {
                                                             private:
                                                              vector< vector< T > > tbl;
class LiChao {
private:
                                                              vector< int > lg;
                                                              T cv(Ta, Tb) {
 int n; vector< Line > nodes;
  inline int lc( int x ) { return 2 * x + 1; }
                                                               return Cmp_()( a, b ) ? a : b;
  inline int rc( int x ) { return 2 * x + 2; }
  void insert( int 1, int r, int id, Line ln ) {
                                                             public:
  int m = (1 + r) >> 1;
                                                              void init( T arr[], int n ) {
  if ( nodes[ id ].id == -1 ) {
                                                               // 0-base
   nodes[ id ] = ln;
                                                               lg.resize( n + 1 );
                                                               lg[0] = -1;
    return:
                                                               for( int i=1 ; i<=n ; ++i ) lg[i] = lg[i>>1] + 1;
   bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
                                                               tbl.resize( lg[n] + 1 );
   if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
                                                               tbl[ 0 ].resize( n );
                                                               copy( arr, arr + n, tbl[ 0 ].begin() );
   atLeft ^= 1; swap( nodes[ id ], ln );
                                                               for ( int i = 1 ; i <= lg[ n ] ; ++ i ) {
  int len = 1 << ( i - 1 ), sz = 1 << i;</pre>
   if ( r - 1 == 1 ) return;
   if ( atLeft ) insert( l, m, lc( id ), ln );
                                                                tbl[ i ].resize( n - sz + 1 );
                                                                for ( int j = 0 ; j \le n - sz ; ++ j
   else insert( m, r, rc( id ), ln );
                                                                 tbl[i][j] = cv(tbl[i-1][j], tbl[i-1][j+len]);
  int query( int 1, int r, int id, int x ) {
                                                               }
   int ret = 0;
                                                              T query( int 1, int r ) {
   if ( nodes[ id ].id != -1 )
                                                               // 0-base [1, r)
    ret = nodes[ id ].at( x );
                                                               int wh = lg[ r - l ], len = 1 << wh;
return cv( tbl[ wh ][ l ], tbl[ wh ][ r - len ] );
   int m = (1 + r) >> 1;
   if ( r - l == 1 ) return ret;
   else if ( x < m )</pre>
                                                              }
                                                            };
    return max( ret, query( 1, m, lc( id ), x ) );
   else
                                                             2.7 Linear Basis
    return max( ret, query( m, r, rc( id ), x ) );
                                                             struct LinearBasis {
public:
                                                             private:
 void build( int n_ ) {
                                                              int n, sz
  n = n_; nodes.clear();
                                                              vector< llu > B;
  nodes.resize( n << 2, Line() );</pre>
                                                              inline llu two( int x ){ return ( ( llu ) 1 ) << x; }</pre>
                                                             public:
 void insert( Line ln ) { insert( 0, n, 0, ln ); }
                                                              void init( int n_ ) {
                                                               n = n_; B.clear(); B.resize( n ); sz = 0;
  int query( int x ) { return query( 0, n, 0, x ); }
} lichao;
                                                              void insert( llu x ) {
2.5 Treap
                                                               // add x into B
namespace Treap{
                                                               for ( int i = n-1; i >= 0; --i ) if( two(i) & x ){
#define sz(x)((x)?((x)-size):0)
                                                                if ( B[ i ] ) x ^= B[ i ];
struct node{
                                                                else {
 int size;
                                                                 B[i] = x; sz++;
                                                                 for ( int j = i - 1 ; j >= 0 ; -- j )
if( B[ j ] && ( two( j ) & B[ i ] ) )
 uint32_t pri;
 node *lc, *rc;
                                                                 B[ i ] ^= B[ j ];
for (int j = i + 1
 node() : size(0), pri(rand()), lc( 0 ), rc( 0 ) {}
                                                                                     ; j < n ; ++ j )
 void pull() {
                                                                  if ( two( i ) & B[ j ] )
   size = 1;
   if ( lc ) size += lc->size;
                                                                   B[ j ] ^= B[ i ];
   if ( rc ) size += rc->size;
                                                                 break;
  }
}:
                                                               }
node* merge( node* L, node* R ) {
 if ( not L or not R ) return L ? L : R;
                                                              inline int size() { return sz; }
                                                              bool check( llu x ) {
  if ( L->pri > R->pri ) {
  L->rc = merge( L->rc, R ); L->pull();
                                                               // is x in span(B) ?
   return L;
                                                               for ( int i = n-1 ; i \ge 0 ; --i ) if( two(i) & x )
```

vector<bool> ap, ins;

```
if( B[ i ] ) x ^= B[ i ];
                                                               void dfs(int u, int f) -
                                                                dfn[u] = low[u] = dfn[f] + 1;
   else return false;
                                                                int ch = 0;
  return true:
                                                                for (auto [v, t]: G[u]) if (v != f) {
  if (not ins[t]) {
llu kth_small(llu k) {
  /** 1-base would always > 0 **/
                                                                  st.push_back(t);
  /** should check it **/
                                                                  ins[t] = true;
  /* if we choose at least one element
   but size(B)(vectors in B)==N(original elements)
                                                                 if (dfn[v]) {
    then we can't get 0 */
                                                                  low[u] = min(low[u], dfn[v]);
  llu ret = 0;
  for ( int i = 0 ; i < n ; ++ i ) if( B[ i ] ) {</pre>
                                                                 } ++ch; dfs(v, u);
  if( k & 1 ) ret ^= B[ i ];
                                                                 low[u] = min(low[u], low[v]);
  k >>= 1;
                                                                 if (low[v] >= dfn[u]) {
                                                                  ap[u] = true;
                                                                  while (true) {
 return ret;
                                                                   int eid = st.back(); st.pop_back();
                                                                   bcc[eid] = ecnt;
} base;
                                                                   if (eid == t) break;
3 Graph
                                                                  ecnt++:
    Euler Circuit
                                                                 }
bool vis[ N ]; size_t la[ K ];
                                                                if (ch == 1 and u == f) ap[u] = false;
void dfs( int u, vector< int >& vec ) {
while ( la[ u ] < G[ u ].size() ) {
  if( vis[ G[ u ][ la[ u ] ].second ] ) {</pre>
                                                              public:
                                                               void init(int n_) {
   ++ la[ u ];
                                                                G.clear(); G.resize(n = n_);
   continue:
                                                                ecnt = 0; ap.assign(n, false);
                                                                low.assign(n, 0); dfn.assign(n, 0);
 int v = G[ u ][ la[ u ] ].first;
 vis[ G[ u ][ la[ u ] ].second ] = true;
                                                               void add_edge(int u, int v) {
 ++ la[ u ]; dfs( v, vec );
                                                                G[u].emplace_back(v, ecnt);
  vec.push_back( v );
                                                                G[v].emplace_back(u, ecnt++);
                                                               void solve() {
                                                                ins.assign(ecnt, false);
3.2 BCC Edge
                                                                bcc.resize(ecnt); ecnt = 0;
class BCC_Bridge {
                                                                for (int i = 0; i < n; ++i)
private:
                                                                 if (not dfn[i]) dfs(i, i);
  int n, ecnt;
 vector<vector<pair<int,int>>> G;
                                                               int get_id(int x) { return bcc[x]; }
  vector<int> dfn, low;
                                                               int count() { return ecnt; }
  vector<bool> bridge;
                                                               bool is_ap(int x) { return ap[x]; }
  void dfs(int u, int f)
                                                             } bcc_ap;
   dfn[u] = low[u] = dfn[f] + 1;
  for (auto [v, t]: G[u]) {
  if (v == f) continue;
                                                             3.4 2-SAT (SCC)
                                                             class TwoSat{
    if (dfn[v]) {
                                                              private:
     low[u] = min(low[u], dfn[v]);
                                                               int n;
     continue;
                                                               vector<vector<int>> rG,G,sccs;
                                                               vector<int> ord,idx;
    dfs(v, u);
low[u] = min(low[u], low[v]);
                                                               vector<bool> vis,result;
                                                               void dfs(int u){
    if (low[v] > dfn[u]) bridge(t) = true;
                                                                vis[u]=true
   }
                                                                for(int v:G[u])
                                                                 if(!vis[v]) dfs(v);
public:
                                                                ord.push_back(u);
  void init(int n_) {
   G.clear(); G.resize(n = n_);
                                                               void rdfs(int u){
  low.assign(n, ecnt = 0);
                                                                vis[u]=false;idx[u]=sccs.size()-1;
   dfn.assign(n, 0);
                                                                sccs.back().push_back(u);
                                                                for(int v:rG[u])
  void add_edge(int u, int v) {
                                                                 if(vis[v])rdfs(v);
   G[u].emplace_back(v, ecnt);
   G[v].emplace_back(u, ecnt++);
                                                              public:
                                                               void init(int n_){
  void solve() {
                                                                n=n_;G.clear();G.resize(n);
  bridge.assign(ecnt, false);
                                                                rG.clear();rG.resize(n)
   for (int i = 0; i < n; ++i)
                                                                sccs.clear();ord.clear();
    if (not dfn[i]) dfs(i, i);
                                                                idx.resize(n);result.resize(n);
  bool is_bridge(int x) { return bridge[x]; }
                                                               void add_edge(int u,int v){
} bcc_bridge;
                                                                G[u].push_back(v);rG[v].push_back(u);
3.3 BCC Vertex
                                                               void orr(int x,int y){
class BCC_AP {
                                                                if ((x^y)==1)return
private:
                                                                add_edge(x^1,y); add_edge(y^1,x);
 int n, ecnt;
 vector<vector<pair<int,int>>> G;
                                                               bool solve(){
 vector<int> bcc, dfn, low, st;
                                                                vis.clear();vis.resize(n);
```

for(int i=0;i<n;++i)</pre>

```
if(not vis[i])dfs(i);
   reverse(ord.begin(),ord.end());
                                                                 void decompose(){
   for (int u:ord){
                                                                  chain_ = 1;
                                                                  predfs( 1, 1 );
    if(!vis[u])continue;
    sccs.push_back(vector<int>());
                                                                  time_{-} = 0;
    rdfs(u);
                                                                  dfschain(1,1);
                                                                 PII get_subtree(int u) { return {tl[ u ],tr[ u ] }; }
   for(int i=0;i<n;i+=2)</pre>
    if(idx[i]==idx[i+1])
                                                                 vector< PII > get_path( int u , int v ){
                                                                  vector< PII > res;
     return false;
                                                                  int g = lca( u, v )
   vector<bool> c(sccs.size());
                                                                  while ( chain[ u ] != chain[ g ] ) {
   for(size_t i=0;i<sccs.size();++i){</pre>
    for(size_t j=0;j<sccs[i].size();++j){
  result[sccs[i][j]]=c[i];</pre>
                                                                   int s = chain_st[ chain[ u ] ];
res.emplace_back( tl[ s ], tl[ u ] + 1 );
     c[idx[sccs[i][j]^1]]=!c[i];
                                                                   u = fa[ s ][ 0 ];
                                                                  res.emplace_back( tl[ g ], tl[ u ] + 1 );
while ( chain[ v ] != chain[ g ] ) {
   return true;
                                                                   int s = chain_st[ chain[ v ] ];
  bool get(int x){return result[x];}
                                                                   res.emplace_back( tl[s], tl[v] + 1);
  inline int get_id(int x){return idx[x];}
                                                                   v = fa[ s ][ 0 ];
  inline int count(){return sccs.size();}
                                                                  res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
                                                                  return res
3.5 Lowbit Decomposition
                                                                  /* res : list of intervals from u to v
class LowbitDecomp{
                                                                   * ( note only nodes work, not edge )
                                                                   * usage :
private:
                                                                   * vector< PII >& path = tree.get_path( u , v )
 int time_, chain_, LOG_N;
                                                                   * for( auto [ 1, r ] : path ) {
* 0-base [ 1, r )
 vector< vector< int > > G, fa;
 vector< int > tl, tr, chain, chain_st;
                                                                   * }
 // chain_ : number of chain
                                                                   */
 // tl, tr[ u ] : subtree interval in the seq. of u
                                                                 }
 // chain_st[ u ] : head of the chain contains u
 // chian[ u ] : chain id of the chain u is on
                                                               } tree;
 void predfs( int u, int f ) {
                                                                     MaxClique
  chain[ u ] = 0;
  for ( int v : G[ u ] ) {
                                                                // contain a self loop u to u, than u won't in clique
   if ( v == f ) continue;
                                                                template < size_t MAXN >
   predfs( v, u );
                                                                class MaxClique{
   if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
                                                                private
    chain[ u ] = chain[ v ];
                                                                 using bits = bitset< MAXN >;
                                                                 bits popped, G[ MAXN ], ans
  if ( not chain[ u ] )
                                                                 size_t deg[ MAXN ], deo[ MAXN ], n;
   chain[ u ] = chain_ ++;
                                                                 void sort_by_degree() {
                                                                  popped.reset();
 void dfschain( int u, int f ) {
                                                                  for ( size_t i = 0 ; i < n ; ++ i )
                                                                  deg[ i ] = G[ i ].count();
for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;</pre>
 fa[ u ][ 0 ] = f;
for ( int i = 1 ; i < LOG_N ; ++ i )
   fa[u][i] = fa[fa[u][i-1]][i-1];
                                                                    for ( size_t j = 0 ; j < n ; ++ j</pre>
  tl[ u ] = time_++
  if ( not chain_st[ chain[ u ] ] )
                                                                       if ( not popped[ j ] and deg[ j ] < mi )</pre>
   chain_st[ chain[ u ] ] = u;
                                                                         mi = deg[id = j]
                                                                    popped[ deo[ i ] = id ] = 1;
  for ( int v : G[ u ] )
   if ( v != f and chain[ v ] == chain[ u ] )
                                                                     for( size_t u = G[ i ]._Find_first() ;
  dfschain( v, u );
for ( int v : G[ u ] )
                                                                     u < n ; u = G[ i ]._Find_next( u ) )
                                                                       -- deg[ u ];
   if ( v != f and chain[ v ] != chain[ u ] )
    dfschain( v, u );
  tr[ u ] = time_;
                                                                 void BK( bits R, bits P, bits X ) {
                                                                  if (R.count()+P.count() <= ans.count()) return;</pre>
 bool anc( int u, int v ) {
  return tl[ u ] <= tl[ v ] and tr[ v ] <= tr[ u ];</pre>
                                                                  if ( not P.count() and not X.count() ) {
                                                                   if ( R.count() > ans.count() ) ans = R;
                                                                   return:
public:
                                                                  }
                                                                   /* greedily chosse max degree as pivot
 int lca( int u, int v ) {
  if ( anc( u, v ) ) return u;
                                                                  bits cur = P | X; size_t pivot = 0, sz = 0;
  for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
                                                                  for ( size_t u = cur._Find_first() ;
   if ( not anc( fa[ u ][ i ], v ) )
u = fa[ u ][ i ];
                                                                   u < n ; u = cur._Find_next( u ) )
if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
                                                                  cur = P & ( ~G[ pivot ] );
  return fa[ u ][ 0 ];
                                                                   */ // or simply choose first
                                                                  bits cur = P & (~G[ ( P | X )._Find_first() ]);
 void init( int n ) {
  fa.assign( ++n, vector< int >( LOG_N ) );
                                                                  for ( size_t u = cur._Find_first()
  for ( LOG_N = 0 ; ( 1 << LOG_N ) < n ; ++ LOG_N );
                                                                   u < n ; u = cur._Find_next( u ) ) {
  G.clear(); G.resize( n );
                                                                   if ( R[ u ] ) continue;
 tl.assign( n, 0 ); tr.assign( n, 0 );
                                                                   R[u] = 1;
                                                                   BK(R, P & G[u], X & G[u]);
R[u] = P[u] = 0, X[u] = 1;
  chain.assig( n, 0 ); chain_st.assign( n, 0 );
 void add_edge( int u , int v ) {
  // 1-base
  G[ u ].push_back( v );
                                                                public:
  G[ v ].push_back( u );
                                                                void init( size_t n_ ) {
```

```
vector<int> r,
                                                                   for (int i = 0; i < n; i++)
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
   G[ i ].reset();
                                                                    if (mask[i]) r.push_back(i);
                                                                   for (int i = 0; i < n; i++)
  ans.reset();
                                                                    d[i] = int((a[i] & mask).count());
void add_edges( int u, bits S ) { G[ u ] = S; }
void add_edge( int u, int v ) {
                                                                   sort(r.begin(), r.end(),
                                                                    [&](int i, int j) { return d[i] > d[j]; });
 G[u][v] = G[v][u] = 1;
                                                                   csort(r, c);
                                                                   dfs(r, c, 1, mask);
                                                                   return ans; // sol[0 ~ ans-1]
int solve() {
  sort_by_degree(); // or simply iota( deo... )
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                } graph;
   deg[ i ] = G[ i ].count();
                                                                 3.8 Virtural Tree
  bits pob, nob = 0; pob.set()
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;
                                                                inline bool cmp(const int &i, const int &j) {
                                                                  return dfn[i] < dfn[j];</pre>
 for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
   size_t v = deo[ i ];
   bits tmp; tmp[ v ] = 1;
                                                                 void build(int vectrices[], int k) {
   BK( tmp, pob & G[ v ], nob & G[ v ] );
pob[ v ] = 0, nob[ v ] = 1;
                                                                  static int stk[MAX_N];
                                                                  sort(vectrices, vectrices + k, cmp);
                                                                  stk[sz++] = 0;
                                                                  for (int i = 0; i < k; ++i) {
  int u = vectrices[i], lca = LCA(u, stk[sz - 1]);
  if (lca == stk[sz - 1]) stk[sz++] = u;</pre>
  return static_cast< int >( ans.count() );
};
     MaxCliqueDyn
                                                                    while (sz >= 2 && dep[stk[sz - 2]] >= dep[lca]) {
                                                                     addEdge(stk[sz - 2], stk[sz - 1]);
constexpr int kN = 150;
struct MaxClique { // Maximum Clique
bitset<kN> a[kN], cs[kN];
int ans, sol[kN], q, cur[kN], d[kN], n;
                                                                    }
                                                                    if (stk[sz - 1] != lca) {
                                                                     addEdge(lca, stk[--sz]);
void init(int _n) {
                                                                     stk[sz++] = lca, vectrices[cnt++] = lca;
 n = _n; for (int i = 0; i < n; i++) a[i].reset();
void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
                                                                    stk[sz++] = u;
void csort(vector<int> &r, vector<int> &c) 
 int mx = 1, km = max(ans - q + 1, 1), t = 0,
                                                                  for (int i = 0; i < sz - 1; ++i)
    m = int(r.size())
  cs[1].reset(); cs[2].reset();
                                                                   addEdge(stk[i], stk[i + 1]);
 for (int i = 0; i < m; i++) {
  int p = r[i], k = 1;</pre>
                                                                 3.9 Centroid Decomposition
   while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
                                                                 struct Centroid {
   cs[k][p] = 1;
                                                                  vector<vector<int64_t>> Dist;
   if (k < km) r[t++] = p;
                                                                  vector<int> Parent, Depth;
                                                                  vector<int64_t> Sub, Sub2;
vector<int> Sz, Sz2;
  c.resize(m);
 if (t) c[t - 1] = 0;
                                                                  Centroid(vector<vector<pair<int, int>>> g) {
  for (int k = km; k <= mx; k++) {</pre>
                                                                   int N = g.size();
   for (int p = int(cs[k]._Find_first());
                                                                   vector<bool> Vis(N);
      p < kN; p = int(cs[k]._Find_next(p))) {</pre>
                                                                   vector<int> sz(N), mx(N);
    r[t] = p; c[t++] = k;
                                                                   vector<int> Path;
                                                                   Dist.resize(N)
  }
                                                                   Parent.resize(N);
                                                                   Depth.resize(N)
                                                                   auto DfsSz = [\&](auto dfs, int x) -> void {
Vis[x] = true; sz[x] = 1; mx[x] = 0;
 void dfs(vector<int> &r, vector<int> &c, int 1,
 bitset<kN> mask) {
                                                                    for (auto [u, w] : g[x]) {
  if (Vis[u]) continue;
 while (!r.empty()) {
   int p = r.back(); r.pop_back();
                                                                     dfs(dfs, u)
   mask[p] = 0;
                                                                     sz[x] += sz[u];
   if (q + c.back() <= ans) return;</pre>
   cur[q++] = p;
                                                                     mx[x] = max(mx[x], sz[u]);
   vector<int> nr, nc;
   bitset<kN> nmask = mask & a[p];
                                                                    Path.push_back(x);
                                                                   };
   for (int i : r)
    if (a[p][i]) nr.push_back(i);
                                                                   auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
   if (!nr.empty()) {
                                                                    -> void {
    if (1 < 4) {
                                                                    Dist[x].push_back(D);Vis[x] = true;
                                                                    for (auto [u, w] : g[x]) {
  if (Vis[u]) continue;
     for (int i : nr)
      d[i] = int((a[i] & nmask).count());
     sort(nr.begin(), nr.end(),
                                                                     dfs(dfs, u, D + w);
      [&](int x, int y) {
  return d[x] > d[y];
                                                                    }
                                                                   };
      });
                                                                   auto Dfs = [&]
                                                                    (auto dfs, int x, int D = 0, int p = -1)->void {
                                                                    Path.clear(); DfsSz(DfsSz, x);
    csort(nr, nc); dfs(nr, nc, l + 1, nmask);
   } else if (q > ans) {
                                                                    int M = Path.size();
                                                                    int C = -1;
    ans = q; copy(cur, cur + q, sol);
                                                                    for (int u : Path) {
   c.pop_back(); q--;
                                                                     if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
 }
                                                                     Vis[u] = false;
int solve(bitset<kN> mask) { // vertex mask
                                                                    DfsDist(DfsDist, C);
```

```
for (int u : Path) Vis[u] = false;
                                                                     if(d[n][i]<inf-eps)</pre>
                                                                      avg=max(avg,(d[n][i]-d[k][i])/(n-k));
   Parent[C] = p; Vis[C] = true;
   Depth[C] = D;
                                                                     else avg=max(avg,inf);
   for (auto [u, w] : g[C]) {
    if (Vis[u]) continue
                                                                    if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
    dfs(dfs, u, D + 1, C);
                                                                   FZ(vst);edgeID.clear();cycle.clear();rho.clear();
                                                                   for (int i=n; !vst[st]; st=prv[i--][st]) {
  Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
                                                                    vst[st]++
  Sz.resize(N); Sz2.resize(N);
                                                                    edgeID.PB(prve[i][st]);
                                                                    rho.PB(st);
 void Mark(int v) {
  int x = v, z = -1;
                                                                   while (vst[st] != 2) {
  for (int i = Depth[v]; i >= 0; --i) {
                                                                    int v = rho.back(); rho.pop_back();
   Sub[x] += Dist[v][i]; Sz[x]++;
                                                                    cycle.PB(v);
   if (z != -1) {
                                                                    vst[v]++;
    Sub2[z] += Dist[v][i];
    Sz2[z]++;
                                                                   reverse(ALL(edgeID));
                                                                   edgeID.resize(SZ(cycle));
   z = x; x = Parent[x];
                                                                   return mmc;
  }
                                                                } mmc;
 int64_t Query(int v) {
                                                                3.12 Mo's Algorithm on Tree
 int64_t res = 0;
  int x = v, z = -1
                                                                int q; vector< int > G[N];
 for (int i = Depth[v]; i >= 0; --i) {
                                                                struct Oue{
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
                                                                  int u, v,
   if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
                                                                 } que[ N ];
                                                                int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
void dfs( int u, int f ) {
  dfn[ u ] = dfn_++; int saved_rbp = stk_;
   z = x; x = Parent[x];
  return res;
                                                                  for ( int v : G[ u ] ) {
 }
                                                                   if ( v == f ) continue;
                                                                   dfs( v, u );
3.10
      Tree Hashing
                                                                  if ( stk_ - saved_rbp < SQRT_N ) continue;
for ( ++ block_ ; stk_ != saved_rbp ; )
  block_id[ stk[ -- stk_ ] ] = block_;</pre>
uint64_t hsah(int u, int f) {
 uint64_t r = 127;
 for (int v : G[ u ]) if (v != f) {
  uint64_t hh = hsah(v, u);
                                                                 stk[ stk_ ++ ] = u;
  r=(r+(hh*hh)%1010101333)%1011820613;
                                                                bool inPath[ N ];
                                                                void Diff( int u ) {
  if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
 return r;
}
                                                                  else { /*add this edge*/ }
3.11 Minimum Mean Cycle
/* minimum mean cycle O(VE) */
                                                                void traverse( int& origin_u, int u ) {
                                                                 for ( int g = lca( origin_u, u )
struct MMC{
                                                                   origin_u != g ; origin_u = parent_of[ origin_u ] )
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
                                                                   Diff( origin_u );
#define V 1021
                                                                  for (int v = u; v != origin_u; v = parent_of[v])
                                                                  Diff( v );
#define inf 1e9
 struct Edge { int v,u; double c; };
int n, m, prv[V][V], prve[V][V], vst[V];
                                                                  origin_u = u;
 Edge e[E];
                                                                void solve() {
                                                                 dfs( 1, 1 );
while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
 vector<int> edgeID, cycle, rho;
 double d[V][V];
 void init( int _n ) { n = _n; m = 0; }
                                                                  sort( que, que + q, [](const Que& x, const Que& y) {
 // WARNING: TYPE matters
                                                                   return tie( block_id[ x.u ], dfn[ x.v ] )
 void add_edge( int vi , int ui ,
                                     double ci )
                                                                       < tie( block_id[ y.u ], dfn[ y.v ] );
 { e[ m ++ ] = { vi , ui , ci }; }
                                                                  int U = 1, V = 1;
for ( int i = 0 ; i < q ; ++ i ) {
 void bellman_ford() {
  for(int i=0; i<n; i++) d[0][i]=0;
  for(int i=0; i<n; i++) {</pre>
                                                                  pass( U, que[ i ].u );
   fill(d[i+1], d[i+1]+n, inf);
for(int j=0; j<m; j++) {
                                                                   pass( V, que[ i ].v );
                                                                   // we could get our answer of que[ i ].id
    int v = e[j].v, u = e[j].u;
                                                                  }
    if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
                                                                }
                                                                /*
     d[i+1][u] = d[i][v]+e[j].c;
     prv[i+1][u] = v;
                                                                Method 2:
     prve[i+1][u] = j;
                                                                dfs u:
                                                                  push u
                                                                  iterate subtree
                                                                 push u
                                                                Let P = LCA(u, v), and St(u) \le St(v)
if (P == u) query[St(u), St(v)]
 double solve(){
  // returns inf if no cycle, mmc otherwise
                                                                else query[Ed(u), St(v)], query[St(P), St(P)]
  double mmc=inf;
  int st = -1;
                                                                3.13 Minimum Steiner Tree
  bellman_ford();
                                                                // Minimum Steiner Tree
  for(int i=0; i<n; i++) {</pre>
   double avg=-inf;
                                                                // 0(V 3^T + V^2 2^T)
   for(int k=0; k<n; k++) {</pre>
                                                                struct SteinerTree{
```

```
#define V 33
#define T 8
                                                                      int x = -1;
#define INF 1023456789
                                                                      for(int i = 1;i <= n;++i)if(i != root && !inc[i]){</pre>
int n , dst[V][V] , dp[1 << T][V] , tdst[V]; void init( int _n ){
                                                                       int j = i, c = 0;
                                                                       while(j!=root && fr[j]!=i && c<=n) ++c, j=fr[j];</pre>
                                                                       if (j == root || c > n) continue;
  for( int i = 0 ; i < n ; i ++ ){</pre>
                                                                       else { x = i; break; }
   for( int j = 0; j < n; j ++)
    dst[ i ][ j ] = INF;
                                                                       if (!~x) {
   dst[ i ][ i ] = 0;
                                                                       for (int i = 1; i <= n; ++i)</pre>
                                                                        if (i != root && !inc[i]) ans += fw[i];
  }
                                                                       return ans;
 void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
                                                                      int y = x;
  dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
                                                                      for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
                                                                      do {
 void shortest_path(){
                                                                       ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true;
  for( int k = 0 ; k < n ; k ++ )
                                                                       } while (y != x);
   for( int i = 0 ; i < n ; i ++ )
                                                                      inc[x] = false;
    for( int j = 0 ; j < n ; j ++ )
dst[ i ][ j ] = min( dst[ i ][ j ],</pre>
                                                                      for (int k = 1; k <= n; ++k) if (vis[k]) {
                                                                       for (int j = 1; j <= n; ++j) if (!vis[j]) {
   if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
   if (g[j][k] < inf && g[j][k]-fw[k] < g[j][x])</pre>
         dst[ i ][ k ] + dst[ k ][ j ] );
                                                                          g[j][x] = g[j][k] - fw[k];
 int solve( const vector<int>& ter ){
  int t = (int)ter.size();
  for( int i = 0 ; i < (1 << t ) ; i ++ )
for( int j = 0 ; j < n ; j ++ )
                                                                      }
                                                                     }
    dp[ i ][ j ] = INF;
                                                                     return ans;
  for( int i = 0 ; i < n ; i ++ )
dp[ 0 ][ i ] = 0;
                                                                    int dfs(int now) {
  for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
                                                                     int r = 1; vis[now] = true;
                                                                     for (int i = 1; i <= n; ++i)
  if (g[now][i] < inf && !vis[i]) r += dfs(i);</pre>
   if( msk == ( msk & (-msk) ) ){
    int who = __lg( msk );
    for( int i = 0 ; i < n ; i ++ )</pre>
                                                                     return r;
     dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
    continue;
                                                                   };
                                                                   3.15
                                                                         Dominator Tree
   for( int i = 0 ; i < n ; i ++ )
                                                                   namespace dominator {
    for( int submsk = ( msk - 1 ) & msk ; submsk ;
          submsk = ( submsk - 1 ) & msk )
                                                                   vector<int> g[maxn], r[maxn], rdom[maxn];
       dp[ msk ][ i ] = min( dp[ msk ][ i ],
                                                                   int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
                dp[ submsk ][ i ]
                                                                   int dom[maxn], val[maxn], rp[maxn], tk;
                dp[ msk ^ submsk ][ i ] );
                                                                   void init(int n) {
   for( int i = 0 ; i < n ; i ++ ){</pre>
                                                                    // vertices are numbered from 0 to n-1
    tdst[ i ] = INF;
for( int j = 0 ; j < n ; j ++ )
tdst[ i ] = min( tdst[ i ],</pre>
                                                                    fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
                                                                    fill(fa, fa + n, -1); fill(val, val + n, -1);
                                                                    fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
                                                                    fill(dom, dom + n, -1); tk = 0;
for (int i = 0; i < n; ++i) {
            dp[ msk ][ j ] + dst[ j ][ i ] );
   for( int i = 0 ; i < n ; i ++ )
                                                                     g[i].clear(); r[i].clear(); rdom[i].clear();
    dp[ msk ][ i ] = tdst[ i ];
                                                                    }
  int ans = INF;
                                                                   void add_edge(int x, int y) { g[x].push_back(y); }
  for( int i = 0 ; i < n ; i ++ )
ans = min( ans , dp[ ( 1 << t ) - 1 ][ i ] );</pre>
                                                                   void dfs(int x) {
                                                                    rev[dfn[x] = tk] = x;
                                                                    fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
  return ans:
                                                                    for (int u : g[x])
} solver;
                                                                     if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
                                                                     r[dfn[u]].push_back(dfn[x]);
      Directed Minimum Spanning Tree
template <typename T> struct DMST {
 T g[maxn][maxn], fw[maxn];
                                                                   void merge(int x, int y) { fa[x] = y; }
 int n, fr[maxn];
                                                                   int find(int x, int c = 0) {
                                                                    if (fa[x] == x) return c ? -1 : x;
 bool vis[maxn], inc[maxn];
                                                                    int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
 void clear() {
  for(int i = 0; i < maxn; ++i) {</pre>
   for(int j = 0; j < maxn; ++j) g[i][j] = inf;</pre>
                                                                    if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
   vis[i] = inc[i] = false;
                                                                    fa[x] = p;
                                                                    return c ? p : val[x];
  }
                                                                   vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
 void addEdge(int u,int v,T w){g[u][v]=min(g[u][v],w);}
 T operator()(int root, int _n) {
  n = n; T ans = 0;
                                                                   // p[i] = -2 if i is unreachable from s
  if (dfs(root) != n) return -1;
                                                                    dfs(s);
                                                                    for (int i = tk - 1; i >= 0; --i)
  while (true) {
                                                                     for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
   for(int i = 1;i <= n;++i) fw[i] = inf, fr[i] = i;</pre>
                                                                     if (i) rdom[sdom[i]].push_back(i);
   for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
    for (int j = 1; j <= n; ++j) +
                                                                     for (int &u : rdom[i]) {
                                                                      int p = find(u);
     if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
                                                                      if (sdom[p] == i) dom[u] = i;
      fw[i] = g[j][i]; fr[i] = j;
                                                                      else dom[u] = p;
```

```
if (i) merge(i, rp[i]);
vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];</pre>
for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
return p;
```

4 Matching & Flow

Kuhn Munkres

```
class KM {
private:
static constexpr lld INF = 1LL << 60;</pre>
vector<lld> hl,hr,slk;
vector<int> fl,fr,pre,qu;
vector<vector<lld>> w;
vector<bool> v1,vr;
 int n, ql, qr;
bool check(int x) {
 if (v1[x] = true, f1[x] != -1)
   return vr[qu[qr++] = f1[x]] = true;
 while (x != -1) swap(x, fr[fl[x] = pre[x]]);
  return false;
 void bfs(int s) {
 fill(slk.begin(), slk.end(), INF);
  fill(vl.begin(), vl.end(), false);
  fill(vr.begin(), vr.end(), false);
  ql = qr = 0;
  qu[qr++] = s;
  vr[s] = true;
  while (true) {
   11d d;
   while (ql < qr) {</pre>
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
     if(!v1[x]&&slk[x]>=(d=h1[x]+hr[y]-w[x][y])){
      if (pre[x] = y, d) slk[x] = d;
      else if (!check(x)) return;
     }
    }
   d = INF;
   for (int x = 0; x < n; ++x)
    if (!vl[x] \&\& d > slk[x]) d = slk[x];
   for (int x = 0; x < n; ++x) {
   if (v1[x]) h1[x] += d;
    else slk[x] -= d;
    if (vr[x]) hr[x] -= d;
   for (int x = 0; x < n; ++x)
    if (!v1[x] && !slk[x] && !check(x)) return;
 }
public:
void init( int n_ ) {
 n = n_; qu.resize(n);
 fl.clear(); fl.resize(n, -1);
 fr.clear(); fr.resize(n, -1);
 hr.clear(); hr.resize(n); hl.resize(n);
 w.clear(); w.resize(n, vector<lld>(n));
 slk.resize(n); pre.resize(n);
 vl.resize(n); vr.resize(n);
 void set_edge( int u, int v, lld x ) {w[u][v] = x;}
11d solve() {
  for (int i = 0; i < n; ++i)</pre>
  hl[i] = *max_element(w[i].begin(), w[i].end());
  for (int i = 0; i < n; ++i) bfs(i);
 11d res = 0;
  for (int i = 0; i < n; ++i) res += w[i][f1[i]];</pre>
} km;
```

4.2 Bipartite Matching

```
class BipartiteMatching{
private:
vector<int> X[N], Y[N];
int fX[N], fY[N], n;
```

```
bitset<N> walked;
 bool dfs(int x)+
  for(auto i:X[x]){
   if(walked[i])continue;
   walked[i]=1
   if(fY[i]==-1||dfs(fY[i])){
    fY[i]=x;fX[x]=i;
    return 1:
  return 0;
public:
 void init(int _n){
  n=_n; walked.reset();
  for(int i=0;i<n;i++){</pre>
   X[i].clear();Y[i].clear();
   fX[i]=fY[i]=-1;
  }
 void add_edge(int x, int y){
 X[x].push_back(y); Y[y].push_back(y);
 int solve(){
  int cnt = 0;
  for(int i=0;i<n;i++){</pre>
   walked.reset();
   if(dfs(i)) cnt++;
  // return how many pair matched
  return cnt;
4.3 General Graph Matching
const int N = 514, E = (2e5) * 2;
struct Graph{
 int to[E],bro[E],head[N],e;
 int lnk[N], vis[N], stp, n;
 void init( int _n ){
  stp = 0; e = 1; n = _n;
  for( int i = 0 ; i <= n ; i ++ )
   head[i] = lnk[i] = vis[i] = 0;
 void add_edge(int u,int v){
  // 1-base
  to[e]=v,bro[e]=head[u],head[u]=e++;
  to[e]=u,bro[e]=head[v],head[v]=e++;
 bool dfs(int x){
  vis[x]=stp;
  for(int i=head[x];i;i=bro[i]){
   int v=to[i]
   if(!lnk[v]){
    lnk[x]=v, lnk[v]=x;
    return true
   }else if(vis[lnk[v]]<stp){</pre>
    int w=lnk[v]
    lnk[x]=v, lnk[v]=x, lnk[w]=0;
    if(dfs(w)) return true
    lnk[w]=v, lnk[v]=w, lnk[x]=0;
  return false;
```

Minimum Weight Matching (Clique version)

```
struct Graph {
 // 0-base (Perfect Match)
 int n, edge[MXN][MXN];
 int match[MXN], dis[MXN], onstk[MXN];
 vector<int> stk;
void init(int _n) {
```

int solve(){

int ans = 0;

return ans;

} } graph;

for(int i=1;i<=n;i++)</pre>

stp++; ans += dfs(i);

if(not lnk[i]){

```
n = _n:
  for (int i=0; i<n; i++)</pre>
   for (int j=0; j<n; j++)</pre>
    edge[i][j] = 0;
void set_edge(int u, int v, int w) {
 edge[u][v] = edge[v][u] = w;
bool SPFA(int u){
 if (onstk[u]) return true;
  stk.PB(u);
  onstk[u] = 1;
 for (int v=0; v<n; v++){
  if (u != v && match[u] != v && !onstk[v]){</pre>
    int m = match[v];
    if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
     dis[m] = dis[u] - edge[v][m] + edge[u][v];
     onstk[v] = 1;
     stk.PB(v);
     if (SPFA(m)) return true;
     stk.pop_back();
     onstk[v] = 0;
  onstk[u] = 0;
  stk.pop_back();
  return false;
int solve() {
  // find a match
  for (int i=0; i<n; i+=2){</pre>
  match[i] = i+1;
  match[i+1] = i;
 while (true){
   int found = 0;
   for (int i=0; i<n; i++)
   dis[i] = onstk[i] = 0;
   for (int i=0; i<n; i++){</pre>
    stk.clear()
    if (!onstk[i] && SPFA(i)){
     found = 1;
     while (SZ(stk)>=2){
      int u = stk.back(); stk.pop_back();
      int v = stk.back(); stk.pop_back();
      match[u] = v;
      match[v] = u;
    }
   if (!found) break;
  int ret = 0;
  for (int i=0; i<n; i++)</pre>
   ret += edge[i][match[i]];
  return ret>>1;
}
} graph;
      Minimum Cost Circulation
struct Edge { int to, cap, rev, cost; };
vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
bool mark[kN];
int NegativeCycle(int n) {
memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
int upd = -1;
for (int i = 0; i <= n; ++i)</pre>
 for (int j = 0; j < n; ++j) {
   int idx = 0;
   for (auto &e : g[j]) {
    if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
     dist[e.to] = dist[j] + e.cost;
```

pv[e.to] = j, ed[e.to] = idx;

while(!mark[upd])mark[upd]=1,upd=pv[upd];

if (i == n) {

return upd;

upd = j;

```
idx++;
   }
  }
 return -1;
int Solve(int n) {
 int rt = -1, ans = 0;
 while ((rt = NegativeCycle(n)) >= 0) {
  memset(mark, false, sizeof(mark));
  vector<pair<int, int>> cyc;
  while (!mark[rt]) {
   cyc.emplace_back(pv[rt], ed[rt]);
   mark[rt] = true;
   rt = pv[rt];
  reverse(cyc.begin(), cyc.end());
  int cap = kInf;
  for (auto &i : cyc)
  auto &e = g[i.first][i.second];
   cap = min(cap, e.cap);
  for (auto &i : cyc) {
   auto &e = g[i.first][i.second];
   e.cap -= cap;
   g[e.to][e.rev].cap += cap;
   ans += e.cost * cap;
 return ans;
}
```

4.6 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}.$

 - 2. For each edge (x,y,l,u), connect $x\to y$ with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect t o s with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T.If $f
 eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the
 - maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f'
 eq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipartite $\operatorname{graph}\left(X,Y\right)$
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect x o y with (cost,cap)=(c,1) if
 - c>0 , otherwise connect $y\to x$ with (cost,cap)=(-c,1) 3. For each edge with c<0 , sum these cost as K , then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost,cap)=0(0, d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =(0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let ${\cal K}$ be the sum of all weights
 - 3. Connect source $s \rightarrow v$, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity
 - 5. For $v \in {\it G}$, connect it with sink $v \to t$ with capacity K+2T- $(\sum_{e \in E(v)} w(e)) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight
 - 2. Connect v
 ightarrow v' with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to \boldsymbol{v}
 - 3. Find the minimum weight perfect matching on G^{\prime} .
- · Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.

- 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
- The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with capacity $c_y. \\$
- 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4.7 Dinic

```
template <typename Cap = int64_t>
class Dinic{
private:
 struct Edge{
  int to, rev;
  Cap cap;
 int n, st, ed;
 vector<vector<Edge>> G;
 vector<int> lv, idx;
 bool BFS(){
  fill(lv.begin(), lv.end(), -1);
  queue<int> bfs;
  bfs.push(st); lv[st] = 0;
  while(!bfs.empty()){
   int u = bfs.front(); bfs.pop();
   for(auto e: G[u]){
    if(e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
    bfs.push(e.to); lv[e.to] = lv[u] + 1;
   }
  }
  return (lv[ed]!=-1);
 Cap DFS(int u, Cap f){
  if(u == ed) return f;
  Cap ret = 0;
  for(int &i = idx[u]; i < (int)G[u].size(); ++i){</pre>
   auto &e = G[u][i];
   if(e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
   Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
   G[e.to][e.rev].cap += nf;
   if(f == 0) return ret;
  if(ret == 0) lv[u] = -1;
  return ret;
public:
 void init(int n_, int st_, int ed_){
  n = n_, st = st_, ed = ed_;
  G.resize(n); lv.resize(n);
  fill(G.begin(), G.end(), vector<Edge>());
void add_edge(int u, int v, Cap c){
G[u].push_back({v, (int)G[v].size(), c});
G[v].push_back({u, ((int)G[u].size())-1, 0});
 Cap max_flow(){
 Cap ret = 0;
  while(BFS()){
   idx.assign(n, 0);
   Cap f = DFS(st, numeric_limits<Cap>::max());
   ret += f;
   if(f == 0) break;
  return ret;
```

4.8 Minimum Cost Maximum Flow

```
class MiniCostMaxiFlow{
  using Cap = int; using Wei = int64_t;
  using PCW = pair<Cap,Wei>;
  static constexpr Cap INF_CAP = 1 << 30;
  static constexpr Wei INF_WEI = 1LL<<60;
  private:</pre>
```

```
struct Edge{
  int to, back;
  Cap cap; Wei wei;
  Edge() {}
  Edge(int a,int b, Cap c, Wei d):
   to(a),back(b),cap(c),wei(d)
  {}
 int ori, edd;
 vector<vector<Edge>> G;
 vector<int> fa, wh;
 vector<bool> inq;
 vector<Wei> dis;
 PCW SPFA(){
  fill(inq.begin(),inq.end(),false);
  fill(dis.begin(), dis.end(), INF_WEI);
  queue<int> qq; qq.push(ori);
  dis[ori]=0;
  while(!qq.empty()){
   int u=qq.front();qq.pop();
   inq[u] = 0;
   for(int i=0;i<SZ(G[u]);++i){</pre>
    Edge e=G[u][i];
    int v=e.to;
    Wei d=e.wei;
    if(e.cap<=0||dis[v]<=dis[u]+d)
     continue:
    dis[v]=dis[u]+d;
    fa[v]=u,wh[v]=i;
    if(inq[v]) continue;
    qq.push(v);
    inq[v]=1;
  if(dis[edd]==INF_WEI) return {-1, -1};
  Cap mw=INF_CAP;
  for(int i=edd;i!=ori;i=fa[i])
   mw=min(mw,G[fa[i]][wh[i]].cap);
  for (int i=edd;i!=ori;i=fa[i]){
   auto &eg=G[fa[i]][wh[i]];
   eg.cap-=mw;
   G[eg.to][eg.back].cap+=mw;
  return {mw,dis[edd]};
public:
 void init(int a,int b,int n){
  ori=a,edd=b;
  G.clear();G.resize(n);
  fa.resize(n);wh.resize(n);
  inq.resize(n); dis.resize(n);
 void add_edge(int st, int ed, Cap c, Wei w){
  G[st].emplace_back(ed,SZ(G[ed]),c,w);
  G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
 PCW solve(){
  /* might modify to
  cc += ret.first * ret.second
  ww += ret.first * ret.second
  */
  Cap cc=0; Wei ww=0;
  while(true){
   PCW ret=SPFA();
   if(ret.first==-1) break;
   cc+=ret.first;
   ww+=ret.second;
  return {cc,ww};
} mcmf;
     Global Min-Cut
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
```

pair<int, int> phase(int n) {
 memset(v, false, sizeof(v));

```
memset(g, 0, sizeof(g));
int s = -1, t = -1;
while (true) {
  int c = -1;
  for (int i = 0; i < n; ++i) {</pre>
   if (del[i] || v[i]) continue;
   if (c == -1 \mid | g[i] > g[c]) c = i;
  if (c == -1) break;
  v[s = t, t = c] = true;
for (int i = 0; i < n; ++i) {
   if (del[i] || v[i]) continue;
   g[i] += w[c][i];
return make_pair(s, t);
int mincut(int n) {
int cut = 1e9;
memset(del, false, sizeof(del));
for (int i = 0; i < n - 1; ++i) {
  int s, t; tie(s, t) = phase(n);
 del[t] = true; cut = min(cut, g[t]);
for (int j = 0; j < n; ++j) {</pre>
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
  }
}
return cut;
```

5 Math

5.1 Prime Table

```
\begin{array}{c} 1002939109, 1020288887, 1028798297, 1038684299, \\ 1041211027, 1051762951, 1058585963, 1063020809, \\ 1147930723, 1172520109, 1183835981, 1187659051, \\ 1241251303, 1247184097, 1255940849, 1272759031, \\ 1287027493, 1288511629, 1294632499, 1312650799, \\ 1868732623, 1884198443, 1884616807, 1885059541, \\ 1909942399, 1914471137, 1923951707, 1925453197, \\ 1979612177, 1980446837, 1989761941, 2007826547, \\ 2008033571, 2011186739, 2039465081, 2039728567, \\ 2093735719, 2116097521, 2123852629, 2140170259, \\ 3148478261, 3153064147, 3176351071, 3187523093, \\ 3196772239, 3201312913, 3203063977, 3204840059, \\ 3210224309, 3213032591, 3217689851, 3218469083, \\ 3219857533, 3231880427, 3235951699, 3273767923, \\ 3276188869, 3277183181, 3282463507, 3285553889, \\ 3319309027, 3327005333, 3327574903, 3341387953, \\ 3373293941, 3380077549, 3380892997, 3381118801 \\ \end{array}
```

5.2 $\lfloor \frac{n}{i} \rfloor$ Enumeration

 $T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor$

5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b==(x/y)*a;
}
```

5.4 Pollard Rho

```
// does not work when n is prime
// return any non-trivial factor
llu pollard_rho(llu n){
static auto f=[](llu x,llu k,llu m){
  return add(k,mul(x,x,m),m);
if (!(n&1)) return 2;
mt19937 rnd(120821011);
while(true){
 llu y=2, yy=y, x=rnd()%n, t=1;
  for(llu sz=2;t==1;sz<<=1) {</pre>
   for(llu i=0;i<sz;++i){</pre>
    if(t!=1)break;
    yy=f(yy,x,n);
    t=gcd(yy>y?yy-y:y-yy,n);
  y=yy;
  if(t!=1&&t!=n) return t;
```

```
5.5
     Pi Count (Linear Sieve)
static constexpr int N = 1000000 + 5;
11d pi[N];
vector<int> primes;
bool sieved[N]
11d cube_root(11d x){
 lld s=cbrt(x-static_cast<long double>(0.1));
 while(s*s*s <= x) ++s;</pre>
 return s-1:
11d square_root(11d x){
 lld s=sqrt(x-static_cast<long double>(0.1));
 while(s*s \ll x) ++s;
 return s-1;
}
void init(){
 primes.reserve(N);
 primes.push_back(1);
for(int i=2;i<N;i++) {</pre>
  if(!sieved[i]) primes.push_back(i);
  pi[i] = !sieved[i] + pi[i-1];
  for(int p: primes) if(p > 1) {
   if(p * i >= N) break;
   sieved[p * i] = true;
   if(p % i == 0) break;
  }
 }
lld phi(lld m, lld n) {
 static constexpr int MM = 80000, NN = 500;
 static lld val[MM][NN];
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
 if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
 return ret;
11d pi_count(11d);
11d P2(11d m, 11d n) {
 11d sm = square_root(m), ret = 0;
 for(lld i = n+1;primes[i]<=sm;i++)</pre>
  ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
 return ret;
11d pi_count(11d m) {
 if(m < N) return pi[m];</pre>
 11d n = pi_count(cube_root(m));
 return phi(m, n) + n - 1 - P2(m, n);
5.6 Range Sieve
const int MAX_SQRT_B = 50000;
const int MAX_L = 200000 + 5;
bool is_prime_small[MAX_SQRT_B];
bool is_prime[MAX_L];
void sieve(lld 1, lld r){
 // [1, r)
 for(lld i=2;i*i<r;i++) is_prime_small[i] = true;</pre>
 for(lld i=1;i<r;i++) is_prime[i-1] = true;</pre>
 if(l==1) is_prime[0] = false;
 for(lld i=2;i*i<r;i++){</pre>
  if(!is_prime_small[i]) continue;
  for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;</pre>
  for(1ld j=std::max(2LL, (1+i-1)/i)*i;j<r;j+=i)</pre>
     is_prime[j-1]=false;
}
5.7 Miller Rabin
bool isprime(llu x){
 static llu magic[]={2,325,9375,28178,\
           450775,9780504,1795265022};
 static auto witn=[](llu a,llu u,llu n,int t)
 ->bool{
  if (!(a = mpow(a%n,u,n)))return 0;
```

while(t--){

1lu a2=mul(a,a,n);

if(a2==1 && a!=1 && a!=n-1)

```
reverse(v.begin() + 1, v.end());
    return 1:
                                                                 for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);
   a = a2:
                                                               VL convolution(const VI &a, const VI &b) {
  return a!=1;
 }:
                                                                // Should be able to handle N <= 10^5, C <= 10^4
 if(x<2)return 0;</pre>
                                                                 int sz = 1;
 if(!(x&1))return x==2;
                                                                 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 llu x1=x-1;int t=0;
                                                                 vector<cplx> v(sz);
 while(!(x1&1))x1>>=1,t++;
                                                                 for (int i = 0; i < sz; ++i) {
                                                                 double re = i < a.size() ? a[i] : 0;
double im = i < b.size() ? b[i] : 0;</pre>
 for(llu m:magic)if(witn(m,x1,x,t))return 0;
return 1;
                                                                  v[i] = cplx(re, im);
5.8 Inverse Element
                                                                 fft(v, sz);
                                                                 for (int i = 0; i <= sz / 2; ++i) {
// x's inverse mod k
                                                                  int j = (sz - i) & (sz - 1);
long long GetInv(long long x, long long k){
                                                                  cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
 // k is prime: euler_(k)=k-1
                                                                    * cplx(0, -0.25);
 return qPow(x, euler_phi(k)-1);
                                                                  if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
// if you need [1, x] (most use: [1, k-1]
void solve(int x, long long k){
                                                                    ].conj()) * cplx(0, -0.25);
                                                                  v[i] = x;
inv[1] = 1;
                                                                 ifft(v, sz);
 for(int i=2;i<x;i++)</pre>
                                                                 VL c(sz):
  inv[i] = ((long long)(k - k/i) * inv[k % i]) % k;
                                                                 for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
                                                                 return c:
5.9 Extended Euler
         \int a^b \mod \varphi(m) + \varphi(m)
                                                                VI convolution_mod(const VI &a, const VI &b, int p) {
                           if (a, m) \neq 1 \land b \geq \varphi(m)
                                                 (\text{mod } m)
         \begin{cases} a^b \mod \varphi(m) \end{cases}
                                                                 int sz = 1:
                           otherwise
                                                                 while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
                                                                 vector<cplx> fa(sz), fb(sz);
5.10 Gauss Elimination
                                                                 for (int i = 0; i < (int)a.size(); ++i)</pre>
void gauss(vector<vector<double>> &d) {
                                                                  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 int n = d.size(), m = d[0].size();
                                                                 for (int i = 0; i < (int)b.size(); ++i)</pre>
 for (int i = 0; i < m; ++i) {
                                                                  fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
  int p = -1;
                                                                 fft(fa, sz), fft(fb, sz);
  for (int j = i; j < n; ++j) {
                                                                 double r = 0.25 / sz;
   if (fabs(d[j][i]) < eps) continue;</pre>
                                                                 cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
   if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
                                                                  int j = (sz - i) & (sz - 1);
  if (p == -1) continue;
                                                                  cplx a1 = (fa[i] + fa[j].conj());
  for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]); for (int j = 0; j < n; ++j) {
                                                                  cplx a2 = (fa[i] - fa[j].conj()) * r2;
cplx b1 = (fb[i] + fb[j].conj()) * r3;
   if (i == j) continue;
                                                                  cplx b2 = (fb[i] - fb[j].conj()) * r4;
                                                                  if (i != j) {
   double z = d[j][i] / d[i][i];
   for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
                                                                   cplx c1 = (fa[j] + fa[i].conj());
                                                                   cplx c2 = (fa[j] - fa[i].conj()) * r2;
                                                                   cplx d1 = (fb[j] + fb[i].conj()) * r3;
                                                                   cplx d2 = (fb[j] - fb[i].conj()) * r4;
                                                                   fa[i] = c1 * d1 + c2 * d2 * r5;
5.11 Fast Fourier Transform
                                                                   fb[i] = c1 * d2 + c2 * d1;
namespace fft {
using VI = vector<int>;
                                                                  fa[j] = a1 * b1 + a2 * b2 * r5;
                                                                  fb[j] = a1 * b2 + a2 * b1;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
                                                                 fft(fa, sz), fft(fb, sz);
                                                                 vector<int> res(sz);
void prefft() {
 generate_n(omega, maxn + 1, [i=0]()mutable{
                                                                 for (int i = 0; i < sz; ++i) {
                                                                  long long a = round(fa[i].re), b = round(fb[i].re),
  auto j = i++;
  return cplx(cos(2*pi*j/maxn), sin(2*pi*j/maxn));
                                                                       c = round(fa[i].im);
                                                                  res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
 });
void fft(vector<cplx> &v, int n) {
                                                                 return res;
 int z = __builtin_ctz(n) - 1;
                                                               }}
 for (int i = 0; i < n; ++i) {
                                                                5.12 Chinese Remainder
  int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^{=(i >> j & 1)<<(z - j);
                                                               1ld crt(lld ans[], lld pri[], int n){
  if (x > i) swap(v[x], v[i]);
                                                                 lld M = 1, ret = 0;
                                                                 for(int i=0;i<n;i++) M *= pri[i];</pre>
 for (int s = 2; s <= n; s <<= 1) {
                                                                 for(int i=0;i<n;i++){</pre>
  int z = s >> 1;
                                                                 lld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
  for (int i = 0; i < n; i += s) {
                                                                  ret += (ans[i]*(M/pri[i])%M * iv)%M;
   for (int k = 0; k < z; ++k) {
                                                                  ret %= M;
    cplx x = v[i + z + k] * omega[maxn / s * k];
    v[i + z + k] = v[i + k] - x;
                                                                 return ret;
    v[i+k] = v[i+k] + x;
                                                                }
                                                                /*
                                                               Another:
                                                               x = a1 \% m1
                                                               x = a2 \% m2
void ifft(vector<cplx> &v, int n) {
                                                               g = gcd(m1, m2)
                                                               assert((a1-a2)%g==0)
fft(v, n);
```

```
[p, q] = exgcd(m2/g, m1/g)
                                                                 VL coef;
return a2+m2*(p*(a1-a2)/g)
                                                                 int n() const { return coef.size(); } // n()>=1
0 <= x < lcm(m1, m2)
                                                                 LL *data() { return coef.data(); }
                                                                 const LL *data() const { return coef.data(); }
                                                                 LL &operator[](size_t i) { return coef[i]; }
      Berlekamp Massey
                                                                 const LL &operator[](size_t i)const{return coef[i];}
                                                                 Poly(initializer_list<LL> a) : coef(a) { }
// x: 1-base, p[]: 0-base
                                                                 explicit Poly(int _n = 1) : coef(_n) { }
template<size_t N>
                                                                 Poly(const LL *arr, int _n) : coef(arr, arr + _n) {}
Poly(const Poly &p, int _n) : coef(_n) {
vector<llf> BM(llf x[N], size_t n){
 size_t f[N]={0},t=0;11f d[N];
                                                                  copy_n(p.data(), min(p.n(), _n), data());
 vector<llf> p[N];
 for(size_t i=1,b=0;i<=n;++i) {</pre>
                                                                 Poly& irev(){return reverse(data(),data()+n()),*this;}
  for(size_t j=0;j<p[t].size();++j)</pre>
                                                                 Poly& isz(int _n) { return coef.resize(_n), *this; }
   d[i]+=x[i-j-1]*p[t][j];
                                                                 Poly& iadd(const Poly &rhs) { // n() == rhs.n()
  if(abs(d[i]-=x[i])<=EPS)continue;</pre>
                                                                  fi(0, n()) if ((coef[i]+=rhs[i]) >= P)coef[i]-=P;
  f[t]=i;if(!t){p[++t].resize(i);continue;}
                                                                  return *this;
  vector<llf> cur(i-f[b]-1);
  11f k=-d[i]/d[f[b]];cur.PB(-k);
                                                                 Poly& imul(LL k) {
  for(size_t j=0;j<p[b].size();j++)</pre>
                                                                  fi(0, n()) coef[i] = coef[i] * k % P;
   cur.PB(p[b][j]*k);
                                                                  return *this:
  if(cur.size()<p[t].size())cur.resize(p[t].size());</pre>
  for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];</pre>
                                                                 Poly Mul(const Poly &rhs) const {
  if(i-f[b]+p[b].size()>=p[t].size()) b=t;
                                                                  const int _n = n2k(n() + rhs.n() - 1);
 p[++t]=cur;
                                                                  Poly X(*this, _n), Y(rhs, _n);
ntt(X.data(), _n), ntt(Y.data(),
fi(0, _n) X[i] = X[i] * Y[i] % P;
return p[t];
                                                                  ntt(X.data(), _n, true);
                                                                  return X.isz(n() + rhs.n() - 1);
5.14 NTT
template <int mod, int G, int maxn>
                                                                 Poly Inv() const { // coef[0] != 0
struct NTT {
                                                                  if (n() == 1) return {ntt.minv(coef[0])};
 static_assert (maxn == (maxn & -maxn));
                                                                  const int _n = n2k(n() * 2);
 int roots[maxn];
                                                                  Poly Xi = Poly(*this, (n() + 1)/2).Inv().isz(_n);
 NTT () {
                                                                  Poly Y(*this, _n);
  int r = modpow(G, (mod - 1) / maxn);
                                                                  ntt(Xi.data(), _n), ntt(Y.data(), _n);
  for (int i = maxn >> 1; i; i >>= 1) {
                                                                  fi(0, _n) {
Xi[i] *= (2 - Xi[i] * Y[i]) % P;
   roots[i] = 1;
   for (int j = 1; j < i; j++)
                                                                   if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
    roots[i + j] = modmul(roots[i + j - 1], r);
   r = modmul(r, r);
                                                                  ntt(Xi.data(), _n, true);
  }
                                                                  return Xi.isz(n());
 // n must be 2^k, and 0 \le F[i] < mod
                                                                 Poly Sqrt() const { // Jacobi(coef[0], P) = 1
 void inplace_ntt(int n, int F[], bool inv = false) {
                                                                  if (n()==1) return {QuadraticResidue(coef[0], P)};
 for (int i = 0, j = 0; i < n; i++) {
  if (i < j) swap(F[i], F[j]);</pre>
                                                                  Poly X = Poly(*this, (n()+1) / 2).Sqrt().isz(n());
                                                                  return X.iadd(Mul(X.Inv()).isz(n())).imul(P/2+1);
   for (int k = n > 1; (j^* = k) < k; k > = 1);
                                                                 pair<Poly, Poly> DivMod(const Poly &rhs) const {
  for (int s = 1; s < n; s *= 2) {
                                                                  // (rhs.)back() != 0
   for (int i = 0; i < n; i += s * 2) {
                                                                  if (n() < rhs.n()) return {{0}, *this};</pre>
    for (int j = 0; j < s; j++) {
                                                                  const int _n = n() - rhs.n() + 1;
     int a = F[i+j];
                                                                  Poly X(rhs); X.irev().isz(_n);
     int b = modmul(F[i+j+s], roots[s+j]);
                                                                  Poly Y(*this); Y.irev().isz(_n);
Poly Q = Y.Mul(X.Inv()).isz(_n).irev();
     F[i+j] = modadd(a, b); // a + b
     F[i+j+s] = modsub(a, b); // a - b
                                                                  X = rhs.Mul(Q), Y = *this;
fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;</pre>
   }
                                                                  return {Q, Y.isz(max(1, rhs.n() - 1))};
  if (inv) {
                                                                 Poly Dx() const {
   int invn = modinv(n);
                                                                  Poly ret(n() - 1);
   for (int i = 0; i < n; i++)
                                                                  fi(0, ret.n()) ret[i] = (i + 1) * coef[i + 1] % P;
    F[i] = modmul(F[i], invn);
                                                                  return ret.isz(max(1, ret.n()));
   reverse(F + 1, F + n);
                                                                 Poly Sx() const {
                                                                  Poly ret(n() + 1);
                                                                  fi(0, n()) ret[i + 1]=ntt.minv(i + 1)*coef[i] % P;
const int P=2013265921, root=31;
                                                                  return ret;
const int MAXN=1<<20;</pre>
NTT<P, root, MAXN> ntt;
                                                                 Poly _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
5.15 Polynomial Operations
                                                                  return Poly(Y.data() + n() - 1, nn);
using VL = vector<LL>;
#define fi(s, n) for (int i=int(s); i<int(n); ++i)
#define Fi(s, n) for (int i=int(n); i>int(s); --i)
                                                                 VL _eval(const VL &x, const auto up)const{
                                                                  const int _n = (int)x.size();
                                                                  if (!_n) return {};
int n2k(int n) {
 int sz = 1; while (sz < n) sz <<= 1;</pre>
                                                                  vector<Poly> down(_n * 2);
                                                                  down[1] = DivMod(up[1]).second;
 return sz;
                                                                  fi(2,_n*2) down[i]=down[i/2].DivMod(up[i]).second;
template<int MAXN, LL P, LL RT> // MAXN = 2^k
                                                                  /* down[1] = Poly(up[1]).irev().isz(n()).Inv().irev()
struct Poly { // coefficients in [0, P)
                                                                     ._tmul(_n, *this);
static NTT<MAXN, P, RT> ntt;
```

```
5.16 FWT
  fi(2, _n * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
    1, down[i / 2]); */
                                                                 /* xor convolution:
  VL y(_n);
                                                                  * x = (x0, x1) , y = (y0, y1)
  fi(0, _n) y[i] = down[_n + i][0];
                                                                  *z = (x0y0 + x1y1, x0y1 + x1y0)
 return y;
                                                                  * =>
}
                                                                  * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
static vector<Poly> _tree1(const VL &x) {
 const int _n = (int)x.size();
                                                                  *z = (1/2) *z'
  vector<Poly> up(_n * 2);
                                                                  * or convolution:
 fi(0, _n) up[_n + i] = \{(x[i] ? P - x[i] : 0), 1\};

Fi(0, _n-1) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
                                                                  * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                                  * and convolution:
 return up;
                                                                 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */const LL MOD = 1e9+7;
VL Eval(const VL&x)const{return _eval(x,_tree1(x));}
                                                                 inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
static Poly Interpolate(const VL &x, const VL &y) {
                                                                  for( int d = 1 ; d < N ; d <<= 1 ) {</pre>
 const int _n = (int)x.size();
                                                                   int d2 = d << 1;
  vector<Poly> up = _tree1(x), down(_n * 2);
                                                                   for( int s = 0; s < N; s += d2)
                                                                    for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
  LL ta = x[ i ] , tb = x[ j ];
  x[ i ] = ta+tb;</pre>
 VL z = up[1].Dx().\_eval(x, up);
 fi(0, _n) z[i] = y[i] * ntt.minv(z[i]) % P;
 fi(0, _n) down[_n + i] = {z[i]};
Fi(0, _n-1) down[i]=down[i * 2].Mul(up[i * 2 + 1])
                                                                     x[ j ] = ta-tb;
   .iadd(down[i * 2 + 1].Mul(up[i * 2]));
                                                                     if( x[ i ] >= MOD ) x[ i ] -= MOD;
  return down[1];
                                                                     if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
                                                                    }
Poly Ln() const \{ // coef[0] == 1 \}
 return Dx().Mul(Inv()).Sx().isz(n());
                                                                  if( inv )
                                                                   for( int i = 0 ; i < N ; i++ ) {
  x[ i ] *= inv( N, MOD );
Poly Exp() const { // coef[0] == 0
 if (n() == 1) return {1};
Poly X = Poly(*this, (n() + 1)/2).Exp().isz(n());
Poly Y = X.Ln(); Y[0] = P - 1;
                                                                    x[ i ] %= MOD;
                                                                 }
  fi(0, n()) if((Y[i] = coef[i] - Y[i]) < 0)Y[i]+=P;
  return X.Mul(Y).isz(n());
                                                                 5.17
                                                                        DiscreteLog
                                                                 11d BSGS(11d P, 11d B, 11d N) {
Poly Pow(const string &K) const {
                                                                  // find B^L = N \mod P
  int nz = 0:
                                                                  unordered_map<lld, int> R;
  while (nz < n() && !coef[nz]) ++nz;</pre>
                                                                  11d sq = (11d)sqrt(P);
 LL nk = 0, nk2 = 0;
                                                                  11d t = 1:
  for (char c : K) {
                                                                  for (int i = 0; i < sq; i++) {
  if (t == N) return i;</pre>
  nk = (nk * 10 + c - '0') % P;
   nk2 = nk2 * 10 + c - '0';
                                                                   if (!R.count(t)) R[t] = i;
   if (nk2 * nz >= n()) return Poly(n());
                                                                   t = (t * B) % P;
   nk2 \ = P - 1;
                                                                  11d f = inverse(t, P);
  if (!nk && !nk2) return Poly({1}, n());
                                                                  for(int i=0;i<=sq+1;i++) {</pre>
 Poly X(data() + nz, n() - nz * nk2);
                                                                   if (R.count(N))
 LL \times 0 = X[0]:
                                                                    return i * sq + R[N];
 return X.imul(ntt.minv(x0)).Ln().imul(nk).Exp()
                                                                   N = (N * f) % P;
   .imul(ntt.mpow(x0, nk2)).irev().isz(n()).irev();
                                                                  return -1;
Poly InvMod(int L) { // (to evaluate linear recursion)
 Poly R{1, 0}; // *this * R mod x^L = 1 (*this[0] ==
                                                                 5.18 FloorSum
  for (int level = 0; (1 << level) < L; ++level) {</pre>
   Poly 0 = R.Mul(Poly(data(), min(2 << level, n())));
                                                                 // @param n `n < 2^32`
                                                                 // @param m `1 <= m < 2^32`
   Poly Q(2 \ll level); Q[0] = 1;
   for (int j = (1 << level); j < (2 << level); ++j)</pre>
                                                                 // @return sum_\{i=0\}^{n-1} floor((ai + b)/m) mod 2^64
    Q[j] = (P - O[j]) \% P;
                                                                 llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
   R = R.Mul(Q).isz(4 << level);
                                                                  llu\ ans = 0;
  }
                                                                  while (true)
                                                                   if (a >= m)
 return R.isz(L):
                                                                    ans += n * (n - 1) / 2 * (a / m); a %= m;
static LL LinearRecursion(const VL&a,const VL&c,LL n){
 // a_n = \sum_{j=0}^{n-j} a_{j}
                                                                   if (b >= m)
  const int k = (int)a.size();
                                                                    ans += n * (b / m); b %= m;
  assert((int)c.size() == k + 1);
 Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
                                                                   llu y_max = a * n + b;
  fi(1, k + 1) C[k - i] = c[i] ? P - c[i] : 0;
                                                                   if (y_max < m) break;</pre>
 C[k] = 1;
                                                                   // y_{max} < m * (n + 1)
  while (n) {
                                                                   // floor(y_max / m) <= n
   if (n % 2) W = W.Mul(M).DivMod(C).second;
                                                                   n = (11u)(y_max / m), b = (11u)(y_max % m);
   n /= 2, M = M.Mul(M).DivMod(C).second;
                                                                   swap(m, a);
  LL ret = 0;
                                                                  return ans;
  fi(0, k) ret = (ret + W[i] * a[i]) % P;
                                                                 11d floor_sum(11d n, 11d m, 11d a, 11d b) {
                                                                  assert(0 <= n && n < (1LL << 32));
assert(1 <= m && m < (1LL << 32));
#undef fi
                                                                  llu ans = 0;
#undef Fi
                                                                  if (a < 0) {
using Poly_t = Poly<131072 * 2, 998244353, 3>;
                                                                   11u \ a2 = (a \% m + m) \% m;
                                                                   ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
template<> decltype(Poly_t::ntt) Poly_t::ntt = {};
```

```
a = a2;
if (b < 0) {
llu b2 = (b % m + m) % m;
 ans -= 1ULL * n * ((b2 - b) / m);
 return ans + floor_sum_unsigned(n, m, a, b);
5.19 Quadratic residue
struct Status{
11 x,y;
11 w;
Status mult(const Status& a,const Status& b,ll mod){
 res.x=(a.x*b.x+a.y*b.y%mod*w)%mod;
 res.y=(a.x*b.y+a.y*b.x)%mod;
 return res;
inline Status qpow(Status _base,11 _pow,11 _mod){
Status res = \{1, 0\};
 while(_pow>0){
 if(_pow&1) res=mult(res,_base,_mod);
  _base=mult(_base,_base,_mod);
  _pow>>=1;
 return res;
inline 11 check(11 x,11 p){
return qpow_mod(x,(p-1)>>1,p);
inline 11 get_root(11 n,11 p){
 if(p==2) return 1
 if(check(n,p)==p-1) return -1;
 ll a;
 while(true){
 a=rand()%p;
  w=((a*a-n)%p+p)%p;
  if(check(w,p)==p-1) break;
 Status res = \{a, 1\}
 res=qpow(res,(p+1)>>1,p);
 return res.x;
5.20 De-Bruijn
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
 if (t > n) {
  if (n \% p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
 } else {
  aux[t] = aux[t - p];
  db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {
   aux[t] = i;
   db(t + 1, t, n, k);
  }
int de_bruijn(int k, int n) {
 // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
 if (k == 1) {
 res[0] = 0;
  return 1;
 for (int i = 0; i < k * n; i++) aux[i] = 0;</pre>
 sz = 0;
 db(1, 1, n, k);
 return sz;
5.21 Simplex Construction
Standard form: maximize \sum_{1 < i < n} c_i x_i such that for all 1 \le j \le m,
\sum_{1 \le i \le n} A_{ji} x_i \le b_j and x_i \ge 0 for all 1 \le i \le n.
```

1. In case of minimization, let $c_i' = -c_i$

```
2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \to \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
3. \sum_{1 \leq i \leq n} A_{ji} x_i = b_j
• \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
• \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j
4. If x_i has no lower bound, replace x_i with x_i - x_i'
```

5.22 Simplex

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d:
vector<int> p, q;
void pivot(int r, int s) {
 double inv = 1.0 / d[r][s];
 for (int i = 0; i < m + 2; ++i)
  for (int j = 0; j < n + 2; ++j)
if (i != r && j != s)
    d[i][j] -= d[r][j] * d[i][s] * inv;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
 d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
 int x = m + z;
 while (true) {
  int s = -1;
  for (int i = 0; i <= n; ++i) {
   if (!z && q[i] == -1) continue
   if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
  if (d[x][s] > -eps) return true;
  int r = -1;
for (int i = 0; i < m; ++i) {
   if (d[i][s] < eps) continue;</pre>
   if (r == -1 ||
    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
  pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
 m = b.size(), n = c.size();
 d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)
for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i) 
p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
 for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];
 q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
 for (int i = 1; i < m; ++i)</pre>
  if (d[i][n + 1] < d[r][n + 1]) r = i;
 if (d[r][n + 1] < -eps) {
  pivot(r, n);
  if (!phase(1) || d[m + 1][n + 1] < -eps)
   return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
   int s = min_element(d[i].begin(), d[i].end() - 1)
        - d[i].begin();
   pivot(i, s);
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)</pre>
  if (p[i] < n) x[p[i]] = d[i][n + 1];
 return x;
}}
```

6 Geometry

6.1 Basic Geometry

```
using coord_t = int;
using Real = double;
using Point = std::complex<coord_t>;
int sgn(coord_t x) {
return (x > 0) - (x < 0);
coord_t dot(Point a, Point b) {
return real(conj(a) * b);
coord_t cross(Point a, Point b) {
return imag(conj(a) * b);
int ori(Point a, Point b, Point c) {
return sgn(cross(b - a, c - a));
bool operator<(const Point &a, const Point &b) {</pre>
return real(a) != real(b)
 ? real(a) < real(b) : imag(a) < imag(b);</pre>
int argCmp(Point a, Point b) {
// -1 / 0 / 1 <-> < / == / > (atan2)
int qa = (imag(a) == 0
  ? (real(a) < 0 ? 3 : 1) : (imag(a) < 0 ? 0 : 2));
int qb = (imag(b) == 0
   ? (real(b) < 0 ? 3 : 1) : (imag(b) < 0 ? 0 : 2));
 if (qa != qb)
 return sgn(qa - qb);
return sgn(cross(b, a));
template <typename V> Real area(const V & pt) {
coord_t ret = 0;
for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
 ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
return ret / 2.0;
```

6.2 Circle Class

```
struct Circle { Point o; Real r; };
vector<Real> intersectAngle(Circle a, Circle b) {
Real d2 = norm(a.o - b.o)
if (norm(A.r - B.r) >= d2)
  if (A.r < B.r)
   return {-PI, PI};
  else
   return {};
 if (norm(A.r + B.r) <= d2) return {};</pre>
Real dis = hypot(A.x - B.x, A.y - B.y);
Real theta = atan^2(B.y - A.y, B.x - A.x);
Real phi = acos((A.r * A.r + d^2 - B.r * B.r) / B.r
   (2 * A.r * dis));
Real L = theta - phi, R = theta + phi;
while (L < -PI) L += PI * 2;
while (R > PI) R -= PI * 2;
return { L, R };
vector<Point> intersectPoint(Circle a, Circle b) {
Real d=o.dis(aa.o);
if (d >= r+aa.r || d <= fabs(r-aa.r)) return {};</pre>
Real dt = (r*r - aa.r*aa.r)/d, d1 = (d+dt)/2;
Point dir = (aa.o-o); dir /= d;
Point pcrs = dir*d1 + o;
dt=sqrt(max(0.0L, r*r - d1*d1)), dir=dir.rot90();
return {pcrs + dir*dt, pcrs - dir*dt};
```

6.3 2D Convex Hull

```
template<typename PT>
vector<PT> buildConvexHull(vector<PT> d) {
  sort(ALL(d), [](const PT& a, const PT& b){
    return tie(a.x, a.y) < tie(b.x, b.y);});
  vector<PT> s(SZ(d)<<1);
  int o = 0;
  for(auto p: d) {
    while(o>=2 && cross(p-s[o-2],s[o-1]-s[o-2])<=0)
    o--;
    s[o++] = p;</pre>
```

```
for(int i=SZ(d)-2, t = o+1;i>=0;i--){
  while(o = t\&cross(d[i] - s[o-2], s[o-1] - s[o-2]) <= 0)
  s[o++] = d[i];
 s.resize(o-1);
 return s;
6.4 3D Convex Hull
// return the faces with pt indexes
int flag[MXN][MXN];
struct Point{
 ld x,y,z;
 Point operator * (const ld &b) const {
  return (Point) {x*b,y*b,z*b};}
 Point operator * (const Point &b) const {
  return(Point){y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
Point ver(Point a, Point b, Point c) {
return (b - a) * (c - a);}
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0
 REP(i,n) REP(j,n) flag[i][j] = 0;
 vector<Face> now;
 now.emplace_back(0,1,2);
 now.emplace_back(2,1,0)
 for (int i=3; i<n; i++){</pre>
  ftop++; vector<Face> next;
REP(j, SZ(now)) {
   Face& f=now[j]; int ff = 0;
   ld d=(pt[i]-pt[f.a]).dot(
     ver(pt[f.a], pt[f.b], pt[f.c]));
   if (d <= 0) next.push_back(f);</pre>
   if (d > 0) ff=ftop;
   else if (d < 0) ff=-ftop
   flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
  REP(j, SZ(now)) {
   Face& f=now[j]
   if (flag[f.a][f.b] > 0 &&
     flag[f.a][f.b] != flag[f.b][f.a])
    next.emplace_back(f.a,f.b,i);
   if (flag[f.b][f.c] > 0 &&
     flag[f.b][f.c] != flag[f.c][f.b])
    next.emplace_back(f.b,f.c,i);
   if (flag[f.c][f.a] > 0 &&
     flag[f.c][f.a] != flag[f.a][f.c])
    next.emplace_back(f.c,f.a,i);
  now=next;
 return now;
6.5 2D Farthest Pair
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {</pre>
while(abs(cross(stk[i+1]-stk[i],
   stk[(pos+1)%n]-stk[i])) >
   abs(cross(stk[i+1]-stk[i],
   stk[pos]-stk[i]))) pos = (pos+1)%n;
 ans = max({ans, dis(stk[i], stk[pos]),
  dis(stk[i+1], stk[pos])});
6.6 2D Closest Pair
struct cmp_y {
bool operator()(const P& p, const P& q) const {
```

return p.y < q.y;</pre>

multiset<P, cmp_y> s;

}):

void solve(P a[], int n) {

sort(a, a + n, [](const P& p, const P& q) {

return tie(p.x, p.y) < tie(q.x, q.y);</pre>

: st(_s), ed(_e), dir(_e - _s) {}

```
llf d = INF; int pt = 0;
 for (int i = 0; i < n; ++i) {
  while (pt < i \text{ and } a[i].x - a[pt].x >= d)
                                                               bool operator<(const Line &lhs, const Line &rhs) {</pre>
   s.erase(s.find(a[pt++]));
                                                                 if (int cmp = argCmp(lhs.dir, rhs.dir))
  auto it = s.lower_bound(P(a[i].x, a[i].y - d));
                                                                   return cmp == -1;
  while (it != s.end() and it->y - a[i].y < d)
                                                                 return ori(lhs.st, lhs.ed, rhs.st) < 0;</pre>
   d = min(d, dis(*(it++), a[i]));
  s.insert(a[i]);
                                                               Point intersect(const Line &A, const Line &B) {
                                                                 Real t = cross(B.st - A.st, B.dir) /
                                                                  cross(A.dir, B.dir);
                                                                 return A.st + t * A.dir;
      kD Closest Pair (3D ver.)
6.7
                                                               }
11f solve(vector<P> v) {
shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,
  unordered_map<lld, int>>> m;
                                                               Real HPI(vector<Line> &lines) {
                                                                 sort(lines.begin(), lines.end());
                                                                 deque<Line> que;
                                                                 deque<Point> pt
 llf d = dis(v[0], v[1]);
                                                                 que.push_back(lines[0]);
auto Idx = [&d] (11f x) -> 11d {
  return round(x * 2 / d) + 0.1; };
                                                                 for (int i = 1; i < (int)lines.size(); i++) {</pre>
                                                                   if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
 auto rebuild_m = [&m, &v, &Idx](int k) {
                                                                    continue
  m.clear();
                                                               #define POP(L, R) \
  for (int i = 0; i < k; ++i)
                                                                   while (pt.size() > 0 \
   m[Idx(v[i].x)][Idx(v[i].y)]
                                                                     && ori(L.st, L.ed, pt.back()) < 0) \
    [Idx(v[i].z)] = i;
                                                                     pt.pop_back(), que.pop_back(); \
 }; rebuild_m(2);
                                                                   while (pt.size() > 0 \
 for (size_t i = 2; i < v.size(); ++i) {
                                                                     && ori(R.st, R.ed, pt.front()) < 0) \
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
                                                                      pt.pop_front(), que.pop_front();
     kz = Idx(v[i].z); bool found = false;
                                                                   POP(lines[i], lines[i])
  for (int dx = -2; dx <= 2; ++dx) {
                                                                   pt.push_back(intersect(que.back(), lines[i]));
   const 11d nx = dx + kx;
                                                                   que.push_back(lines[i]);
   if (m.find(nx) == m.end()) continue;
   auto& mm = m[nx];
                                                                 POP(que.front(), que.back())
   for (int dy = -2; dy <= 2; ++dy) {
                                                                 if (que.size() <= 1 ||</pre>
    const 11d ny = dy + ky;
                                                                   argCmp(que.front().dir, que.back().dir) == 0)
    if (mm.find(ny) == mm.end()) continue;
    auto& mmm = mm[ny];
                                                                 pt.push_back(intersect(que.front(), que.back()));
    for (int dz = -2; dz <= 2; ++dz) {
                                                                 return area(pt);
     const 11d nz = dz + kz;
     if (mmm.find(nz) == mmm.end()) continue;
     const int p = mmm[nz];
                                                               6.10 Minkowski sum
     if (dis(v[p], v[i]) < d) {
  d = dis(v[p], v[i]);</pre>
                                                               vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
      found = true;
                                                                hull(A), hull(B);
                                                                vector<pll> C(1, A[0] + B[0]), s1, s2;
for(int i = 0; i < SZ(A); ++i)</pre>
                                                                 s1.pb(A[(i + 1) % SZ(A)] - A[i]);
                                                                for(int i = 0; i < SZ(B); i++)
s2.pb(B[(i + 1) % SZ(B)] - B[i]);</pre>
  if (found) rebuild_m(i + 1);
  else m[kx][ky][kz] = i;
                                                                for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
                                                                 if (p2 >= SZ(B)
                                                                   | | (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
 return d;
                                                                  C.pb(C.back() + s1[p1++]);
                                                                 else
6.8 Simulated Annealing
                                                                  C.pb(C.back() + s2[p2++]);
                                                                return hull(C), C;
11f anneal() {
 mt19937 rnd_engine( seed );
 uniform_real_distribution< llf > rnd( 0, 1 );
                                                               6.11 intersection of line and circle
 const llf dT = 0.001;
 // Argument p
                                                               vector<pdd> line_interCircle(const pdd &p1,
 11f S_cur = calc( p ), S_best = S_cur;
                                                                   const pdd &p2,const pdd &c,const double r){
 for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
                                                                pdd ft=foot(p1,p2,c),vec=p2-p1;
                                                                double dis=abs(c-ft);
  // Modify p to p_prime
  const 11f S_prime = calc( p_prime );
                                                                if(fabs(dis-r)<eps) return vector<pdd>{ft};
  const llf delta_c = S_prime - S_cur
                                                                if(dis>r) return {};
  llf prob = min( ( llf ) 1, exp(_-delta_c / T ) );
                                                                vec=vec*sqrt(r*r-dis*dis)/abs(vec);
  if ( rnd( rnd_engine ) <= prob )</pre>
                                                                return vector<pdd>{ft+vec,ft-vec};
   S_{cur} = S_{prime}, p = p_{prime};
  if ( S_prime < S_best ) // find min</pre>
                                                               6.12 intersection of polygon and circle
   S_best = S_prime, p_best = p_prime;
                                                               // Divides into multiple triangle, and sum up
 return S_best;
                                                               // test by HDU2892
                                                               const double PI=acos(-1);
                                                               double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
6.9 Half Plane Intersection
// NOTE: Point is complex<Real>
                                                                if(abs(pb)<eps) return 0;</pre>
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
                                                                double S, h, theta;
struct Line {
                                                                double a=abs(pb), b=abs(pa), c=abs(pb-pa);
  Point st, ed;
                                                                double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  Point dir:
                                                                double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  Line (Point _s, Point _e)
                                                                if(a > r){
```

S = (C/2)*r*r;

```
h = a*b*sin(C)/c;
if (h < r && B < PI/2)
S -= (acos(h/r)*r*r - h*sqrt(r*r-h*h));
}
else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
    S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
}
else S = .5*sin(C)*a*b;
return S;
}
double area_poly_circle(const vector<pdd> poly,
    const pdd &0, const double r){
    double S=0;
    for(int i=0;i<SZ(poly);++i)
    S+=_area(poly[i]-0,poly[(i+1)%SZ(poly)]-0,r)
        *ori(0,poly[i],poly[(i+1)%SZ(poly)]);
return fabs(S);
}</pre>
```

6.13 intersection of two circle

```
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
  pdd o1 = a.0, o2 = b.0;
  double r1 = a.R, r2 = b.R, d2 = abs2(o1 - o2),
      d = sqrt(d2);
  if(d < max(r1, r2) - min(r1, r2) || d > r1 + r2)
  return 0;
  pdd u = (o1 + o2) * 0.5
  + (o1 - o2) * ((r2 * r2 - r1 * r1) / (2 * d2));
  double A = sqrt((r1 + r2 + d) * (r1 - r2 + d)
      * (r1 + r2 - d) * (-r1 + r2 + d));
  pdd v = pdd(o1.Y - o2.Y, -o1.X + o2.X) * A
  / (2 * d2);
  p1 = u + v, p2 = u - v;
  return 1;
}
```

6.14 tangent line of two circle

```
vector<Line> go(const Cir& c1,
  const Cir& c2, int sign1){
 // sign1 = 1 for outer tang, -1 for inter tang
vector<Line> ret;
 double d_sq = norm2( c1.0 - c2.0 );
if( d_sq < eps ) return ret;</pre>
double d = sqrt( d_sq );
Pt v = (c2.0 - c1.0) / d;
double c = ( c1.R - sign1 * c2.R ) / d;
 if( c * c > 1 ) return ret;
double h = sqrt( max( 0.0 , 1.0 - c * c ) );
for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
 Pt n = \{ v.X * c - sign2 * h * v.Y ,
  v.Y * c + sign2 * h * v.X };
 Pt p1 = c1.0 + n * c1.R;
 Pt p2 = c2.0 + n * (c2.R * sign1);
 if(fabs(p1.X - p2.X) < eps and
   fabs( p1.Y - p2.Y ) < eps )
   p2 = p1 + perp(c2.0 - c1.0);
 ret.push_back( { p1 , p2 } );
return ret:
```

6.15 Minimum Covering Circle

```
template<typename P>
Circle getCircum(const P &a, const P &b, const P &c){
Real a1 = a.x-b.x, b1 = a.y-b.y;
Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
Real a2 = a.x-c.x, b2 = a.y-c.y;
Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
Circle cc;
cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
return cc;
template<typename P>
Circle MinCircleCover(const vector<P>& pts){
random_shuffle(pts.begin(), pts.end());
Circle c = { pts[0], 0 };
for(int i=0;i<(int)pts.size();i++){</pre>
```

```
if (dist(pts[i], c.o) <= c.r) continue;</pre>
  c = { pts[i], 0 };
  for (int j = 0; j < i; j++) {
   if(dist(pts[j], c.o) <= c.r) continue;
c.o = (pts[i] + pts[j]) / 2;</pre>
   c.r = dist(pts[i], c.o);
   for (int k = 0; k < j; k++) {
    if (dist(pts[k], c.o) <= c.r) continue;</pre>
    c = getCircum(pts[i], pts[j], pts[k]);
  }
 return c;
6.16 KDTree (Nearest Point)
const int MXN = 100005;
struct KDTree {
 struct Node {
  int x,y,x1,y1,x2,y2;
 int id,f;
Node *L, *R;
 } tree[MXN], *root;
 int n
 LL dis2(int x1, int y1, int x2, int y2) {
 LL dx = x1-x2, dy = y1-y2;
  return dx*dx+dy*dy;
 static bool cmpx(Node& a, Node& b){return a.x<b.x;}
static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 void init(vector<pair<int,int>> ip) {
  n = ip.size();
  for (int i=0; i<n; i++) {
   tree[i].id = i;
   tree[i].x = ip[i].first;
   tree[i].y = ip[i].second;
  root = build_tree(0, n-1, 0);
 Node* build_tree(int L, int R, int d) {
  if (L>R) return nullptr;
  int M = (L+R)/2; tree[M].f = d%2;
  nth_element(tree+L,tree+M,tree+R+1,d%2?cmpy:cmpx);
  tree[M].x1 = tree[M].x2 = tree[M].x;
  tree[M].y1 = tree[M].y2 = tree[M].y;
  tree[M].L = build\_tree(L, M-1, d+1);
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
   tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
   tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
  tree[M].R = build_tree(M+1, R, d+1);
  if (tree[M].R) {
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
   tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
   tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
  return tree+M;
 int touch(Node* r, int x, int y, LL d2){
 LL dis = sqrt(d2)+1;
  if (x<r->x1-dis || x>r->x2+dis ||
    y<r->y1-dis || y>r->y2+dis)
   return 0;
  return 1;
 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
  if (!r || !touch(r, x, y, md2)) return;
  LL d2 = dis2(r->x, r->y, x, y);
  if (d2 < md2 || (d2 == md2 && mID < r->id)) {
   mID = r->id;
   md2 = d2;
  // search order depends on split dim
  if ((r->f == 0 && x < r->x) ||
    (r->f == 1 \&\& y < r->y)) {
   nearest(r\rightarrow L, x, y, mID, md2);
   nearest(r->R, x, y, mID, md2);
  } else {
   nearest(r->R, x, y, mID, md2);
```

sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);

pre(sa, c, n, z);

```
nearest(r->L, x, y, mID, md2);
                                                                                         i >= 0; --i
                                                                  for (int i = nn - 1;
                                                                   sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
                                                                  induce(sa, c, s, t, n, z);
 int query(int x, int y) {
 int id = 1029384756;
                                                                 void build(const string &s) {
 LL d2 = 102938475612345678LL;
                                                                  for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];
                                                                  _s[(int)s.size()] = 0; // s shouldn't contain 0
 nearest(root, x, y, id, d2);
                                                                  sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
  return id;
                                                                  for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];
for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
} tree;
                                                                  int ind = 0; hi[0] = 0;
                                                                  for (int i = 0; i < (int)s.size(); ++i) {</pre>
     Stringology
                                                                   if (!rev[i]) {
7.1 Hash
                                                                    ind = 0:
                                                                    continue;
class Hash {
 private:
                                                                   while (i + ind < (int)s.size() && \</pre>
  static constexpr int P = 127, Q = 1051762951;
                                                                    s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  vector<int> h, p;
                                                                   hi[rev[i]] = ind ? ind-- : 0;
 public:
  void init(const string &s){
   h.assign(s.size()+1, 0); p.resize(s.size()+1);
   for (size_t i = 0; i < s.size(); ++i)</pre>
                                                                 7.3 Aho-Corasick Algorithm
    h[i + 1] = add(mul(h[i], P), s[i]);
                                                                 class AhoCorasick{
   generate(p.begin(), p.end(),[x=1,y=1,this]()
                                                                  private:
     mutable{y=x;x=mul(x,P);return y;});
                                                                   static constexpr int Z = 26;
                                                                   struct node{
  int query(int 1, int r){ // 1-base (1, r]
                                                                    node *nxt[ Z ], *fail;
   return sub(h[r], mul(h[1], p[r-1]));}
                                                                    vector< int > data;
                                                                    node(): fail( nullptr ) {
                                                                     memset( nxt, 0, sizeof( nxt ) );
7.2 Suffix Array
                                                                     data.clear();
namespace sfxarray {
                                                                   } *rt;
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
                                                                   inline int Idx( char c ) { return c - 'a'; }
int x[maxn], p[maxn], q[maxn * 2];
                                                                   void init() { rt = new node(); }
// sa[i]: sa[i]-th suffix is the \
                                                                   void add( const string& s, int d ) {
// i-th lexigraphically smallest suffix.
                                                                    node* cur = rt;
// hi[i]: longest common prefix \
                                                                    for ( auto c : s )
// of suffix sa[i] and suffix sa[i - 1].
                                                                     if ( not cur->nxt[ Idx( c ) ] )
void pre(int *sa, int *c, int n, int z) {
                                                                      cur->nxt[ Idx( c ) ] = new node();
                                                                     cur = cur->nxt[ Idx( c ) ];
 memset(sa, 0, sizeof(int) * n);
 memcpy(x, c, sizeof(int) * z);
                                                                    cur->data.push_back( d );
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
if (sa[i] && !t[sa[i] - 1])
                                                                   void compile() {
                                                                    vector< node* > bfs;
                                                                     size_t ptr = 0;
   sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
                                                                    for ( int i = 0 ; i < Z ; ++ i ) {
 memcpy(x, c, sizeof(int) * z);
                                                                     if ( not rt->nxt[ i ] ) {
 for (int i = n - 1; i >= 0; --i)
                                                                      // uncomment 2 lines to make it DFA
  if (sa[i] && t[sa[i] - 1])
                                                                      // rt->nxt[i] = rt;
   sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
                                                                      continue;
void sais(int *s, int *sa, int *p, int *q,
                                                                     rt->nxt[ i ]->fail = rt;
 bool *t, int *c, int n, int z) {
                                                                     bfs.push_back( rt->nxt[ i ] );
 bool uniq = t[n - 1] = true;
 int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
                                                                    while ( ptr < bfs.size() ) {</pre>
                                                                     node* u = bfs[ ptr ++ ];
for ( int i = 0 ; i < Z ; ++ i ) {
  if ( not u->nxt[ i ] ) {
 memset(c, 0, sizeof(int) * z);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
 if (uniq) {
                                                                        // u->nxt[i] = u->fail->nxt[i];
  for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
                                                                        continue;
  return;
                                                                      node* u_f = u->fail;
 for (int i = n - 2; i \ge 0; --i)
                                                                      while ( u_f ) {
 t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
                                                                        if ( not u_f->nxt[ i ] ) {
 pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i)</pre>
                                                                        u_f = u_f->fail; continue;
  if (t[i] && !t[i - 1])
                                                                        u->nxt[ i ]->fail = u_f->nxt[ i ];
   sa[--x[s[i]]] = p[q[i] = nn++] = i;
                                                                        break;
 induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) {
  if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {</pre>
                                                                       if ( not u_f ) u->nxt[ i ]->fail = rt;
                                                                      bfs.push_back( u->nxt[ i ] );
  bool neq = last < 0 ||
   memcmp(s + sa[i], s + last,
(p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
  ns[q[last = sa[i]]] = nmxz += neq;
                                                                   void match( const string& s, vector< int >& ret ) {
                                                                    node* u = rt;
```

for (auto c : s) {

while (u != rt and not u->nxt[Idx(c)])

```
u = u->fail;
                                                               f[i] = k;
    u = u - nxt[Idx(c)];
    if ( not u ) u = rt;
                                                              return f;
    node* tmp = u;
    while ( tmp != rt ) {
                                                             vector<int> search(const string &s, const string &t) {
     for ( auto d : tmp->data )
                                                              // return 0-indexed occurrence of t in s
                                                              vector<int> f = kmp(t), r;
for (int i = 0, k = 0; i < (int)s.size(); ++i)</pre>
      ret.push_back( d );
     tmp = tmp->fail;
                                                               while(k > 0 && (k==(int)t.size() \mid \mid s[i]!=t[k]))
                                                                k = f[k - 1]
  }
                                                               if (s[i] == t[k]) ++k;
                                                               if (k == (int)t.size()) r.push_back(i-t.size()+1);
} ac;
7.4 Suffix Automaton
                                                              return res;
                                                             }
struct Node{
 Node *green, *edge[26];
                                                             7.6 Z value
 int max_len;
 Node(const int _max_len)
                                                             char s[MAXN];
  : green(NULL), max_len(_max_len){
                                                             int len,z[MAXN];
  memset(edge,0,sizeof(edge));
                                                             void Z_value() {
                                                              int i,j,left,right;
} *ROOT, *LAST;
                                                              z[left=right=0]=len;
void Extend(const int c) {
                                                              for(i=1;i<len;i++)</pre>
 Node *cursor = LAST;
                                                               j=max(min(z[i-left],right-i),0);
 LAST = new Node((LAST->max_len) + 1);
                                                               for(;i+j<len&&s[i+j]==s[j];j++);
 for(;cursor&&!cursor->edge[c]; cursor=cursor->green)
                                                               if(i+(z[i]=j)>right)right=i+z[left=i];
  cursor->edge[c] = LAST;
 if (!cursor)
                                                             }
 LAST->green = ROOT;
 else {
                                                             7.7 Manacher
  Node *potential_green = cursor->edge[c];
                                                             int z[maxn];
  if((potential_green->max_len)==(cursor->max_len+1))
                                                             int manacher(const string& s) {
  string t = ".";
   LAST->green = potential_green;
  else {
                                                              for(char c: s) t += c, t += '.';
//assert(potential_green->max_len>(cursor->max_len+1));
                                                              int 1 = 0, r = 0, ans = 0;
   Node *wish = new Node((cursor->max_len) + 1);
                                                              for (int i = 1; i < t.length(); ++i) {
z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
   for(;cursor && cursor->edge[c]==potential_green;
      cursor = cursor->green)
                                                               while (i - z[i] >= 0 \&\& i + z[i] < t.length()) {
    cursor->edge[c] = wish;
                                                                if(t[i - z[i]] == t[i + z[i]]) ++z[i];
   for (int i = 0; i < 26; i++)
                                                                else break;
    wish->edge[i] = potential_green->edge[i];
   wish->green = potential_green->green;
                                                               if (i + z[i] > r) r = i + z[i], l = i;
   potential_green->green = wish;
   LAST->green = wish;
                                                              for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
                                                              return ans;
char S[10000001], A[10000001];
                                                                  Lexico Smallest Rotation
                                                             string mcp(string s){
int main(){
 scanf("%d%s", &N, S);
ROOT = LAST = new Node(0);
                                                              int n = s.length();
                                                              s += s;
                                                              int i=0, j=1;
 for (int i = 0; S[i]; i++)
                                                              while (i<n && j<n){</pre>
  Extend(S[i] - 'a');
                                                               int k = 0;
 while (N--){
  scanf("%s", A);
                                                               while (k < n \&\& s[i+k] == s[j+k]) k++;
                                                               if (s[i+k] <= s[j+k]) j += k+1;</pre>
  Node *cursor = ROOT;
                                                               else i += k+1;
  bool ans = true;
                                                               if (i == j) j++;
  for (int i = 0; A[i]; i++){
   cursor = cursor->edge[A[i] - 'a'];
                                                              int ans = i < n ? i : j;</pre>
   if (!cursor) {
                                                              return s.substr(ans, n);
    ans = false;
    break;
   }
                                                             7.9 BWT
  puts(ans ? "Yes" : "No");
                                                             struct BurrowsWheeler{
                                                             #define SIGMA 26
 return 0;
                                                             #define BASE
                                                              vector<int> v[ SIGMA ];
}
                                                              void BWT(char* ori, char* res){
7.5 KMP
                                                               // make ori -> ori + ori
                                                               // then build suffix array
vector<int> kmp(const string &s) {
 vector<int> f(s.size(), 0);
 /* f[i] = length of the longest prefix
                                                              void iBWT(char* ori, char* res){
                                                               for( int i = 0 ; i < SIGMA ; i ++ )</pre>
   (excluding s[0:i]) such that it coincides
   with the suffix of s[0:i] of the same length */
                                                                v[ i ].clear();
 /* i + 1 - f[i] is the length of the
                                                               int len = strlen( ori );
                                                               for( int i = 0 ; i < len ; i ++ )</pre>
   smallest recurring period of s[0:i] */
                                                                v[`ori[i] - BASE ].push_back( i );
 int k = 0:
 for (int i = 1; i < (int)s.size(); ++i) {</pre>
                                                               vector<int> a;
  while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
                                                               for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
                                                                for( auto j : v[ i ] ){
  if (s[i] == s[k]) ++k;
```

```
a.push_back( j );
  ori[ ptr ++ ] = BASE + i;
}
for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];
  ptr = a[ ptr ];
}
res[ len ] = 0;
}
bwt;</pre>
```

7.10 Palindromic Tree

```
struct palindromic_tree{
struct node{
 int next[26],f,len;
 int cnt, num, st, ed;
 node(int l=0):f(0),len(1),cnt(0),num(0) {
  memset(next, 0, sizeof(next)); }
vector<node> st;
vector<char> s;
int last,n;
void init(){
 st.clear();s.clear();last=1; n=0;
 st.push_back(0);st.push_back(-1);
 st[0].f=1;s.push_back(-1); }
int getFail(int x){
 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
 return x;}
void add(int c){
 s.push_back(c-='a'); ++n;
 int cur=getFail(last);
 if(!st[cur].next[c]){
  int now=st.size();
  st.push_back(st[cur].len+2);
  st[now].f=st[getFail(st[cur].f)].next[c];
  st[cur].next[c]=now;
  st[now].num=st[st[now].f].num+1;
 last=st[cur].next[c];
 ++st[last].cnt;}
int size(){ return st.size()-2;}
} pt;
int main() {
string s; cin >> s; pt.init();
for (int i=0; i<SZ(s); i++) {</pre>
 int prvsz = pt.size(); pt.add(s[i]);
 if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [1,r]: s.substr(1, r-1+1)
 }
return 0;
```

8 Misc

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\dots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \le k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

holds for all 1 < k < n.

8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W\subseteq X, |W|\leq |N_G(W)|$

8.1.7 Euler's planar graph formula

```
V - E + F = C + 1, E < 3V - 6(?)
```

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Lucas's theorem

```
\binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p}, where m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0, and n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0.
```

8.2 MaximumEmptyRect

```
int max_empty_rect(int n, int m, bool blocked[N][N]) {
 static int mxu[2][N], me=0, he=1, ans=0;
 for (int i=0;i<m;i++) mxu[he][i]=0;</pre>
 for (int i=0;i<n;i++) {</pre>
  stack<PII, vector<PII>> stk;
  for (int j=0;j<m;++j) {</pre>
   if (blocked[i][j]) mxu[me][j]=0;
   else mxu[me][j]=mxu[he][j]+1;
   int la = j;
   while (!stk.empty()&&stk.top().FF>mxu[me][j]) {
    int x1 = i - stk.top().FF, x2 = i;
    int y1 = stk.top().SS, y2 = j;
    la = stk.top().SS; stk.pop()
    ans=max(ans,(x2-x1)*(y2-y1));
   if (stk.empty()||stk.top().FF<mxu[me][j])</pre>
    stk.push({mxu[me][j],la});
  while (!stk.empty()) {
   int x1 = i - stk.top().FF, x2 = i;
   int y1 = stk.top().SS-1, y2 = m-1
   stk.pop(); ans=max(ans,(x2-x1)*(y2-y1));
  swap(me, he);
 return ans;
```

8.3 DP-opt Condition

8.3.1 totally monotone (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

8.3.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

8.4 Convex 1D/1D DP

```
struct segment {
  int i, 1, r;
  segment() {}
  segment(int a, int b, int c): i(a), l(b), r(c) {}
};
  inline lld f(int l, int r){return dp[l] + w(l+1, r);}
  void solve() {
   dp[0] = 0;
   deque<segment> dq; dq.push_back(segment(0, 1, n));
   for (int i = 1; i <= n; ++i) {
    dp[i] = f(dq.front().i, i);
    while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
```

```
dq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
                                                             // died at kth
  while (dq.size() &&
                                                            int kth(int n, int m, int k){
   f(i, dq.back().1) < f(dq.back().i, dq.back().1))
                                                              if (m == 1) return n-1;
   dq.pop_back();
                                                              for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  if (dq.size())
                                                              return k;
   int d = 1 << 20, c = dq.back().1;</pre>
   while (d >>= 1) if (c + d <= dq.back().r)</pre>
                                                             8.7 Cactus Matching
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
                                                             vector<int> init_g[maxn],g[maxn*2];
                                                             int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
  if (seg.1 <= n) dq.push_back(seg);</pre>
                                                             void tarjan(int u){
                                                              dfn[u]=low[u]=++dfs_idx;
                                                              for(int i=0;i<(int)init_g[u].size();i++){</pre>
                                                               int v=init_g[u][i];
8.5 ConvexHull Optimization
                                                               if(v==par[u]) continue;
inline 1ld DivCeil(1ld n, 1ld d) { // ceil(n/d)
                                                               if(!dfn[v]){
return n / d + (((n < 0) != (d > 0)) && (n % d));
                                                                par[v]=u;
                                                                tarjan(v);
struct Line {
                                                                low[u]=min(low[u],low[v]);
static bool flag;
                                                                if(dfn[u]<low[v]){</pre>
lld a, b, l, r; // y=ax+b in [l, r)
                                                                 g[u].push_back(v);
lld operator()(lld x) const { return a * x + b; }
bool operator<(const Line& i) const {</pre>
                                                                 g[v].push_back(u);
  return flag ? tie(a, b) < tie(i.a, i.b) : 1 < i.l;</pre>
                                                               }else{
                                                                low[u]=min(low[u],dfn[v]);
11d operator&(const Line& i) const {
                                                                if(dfn[v]<dfn[u]){</pre>
 return DivCeil(b - i.b, i.a - a);
                                                                 int temp_v=u;
                                                                 bcc_id++
                                                                 while(temp_v!=v){
bool Line::flag = true;
                                                                  g[bcc_id+n].push_back(temp_v);
class ConvexHullMax {
                                                                  g[temp_v].push_back(bcc_id+n);
set<Line> L;
                                                                  temp_v=par[temp_v];
public:
ConvexHullMax() { Line::flag = true; }
void InsertLine(lld a, lld b) { // add y = ax + b
                                                                 g[bcc_id+n].push_back(v);
                                                                 g[v].push_back(bcc_id+n);
 Line now = \{a, b, -INF, INF\};
                                                                 reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
  if (L.empty()) {
  L.insert(now);
  return;
                                                             int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
 Line::flag = true;
  auto it = L.lower_bound(now);
                                                             void dfs(int u,int fa){
  auto prv = it == L.begin() ? it : prev(it);
                                                              if(u<=n){
  if (it != L.end() && ((it != L.begin() &&
                                                               for(int i=0;i<(int)g[u].size();i++){</pre>
   (*it)(it->1) >= now(it->1) &&
                                                                int v=g[u][i];
   (*prv)(prv->r - 1) >= now(prv->r - 1)) ||
                                                                if(v==fa) continue;
   dfs(v,u):
  if (it != L.begin()) {
                                                                memset(tp,0x8f,sizeof tp);
   while (prv != L.begin() &&
                                                                if(v<=n){
    (*prv)(prv->1) <= now(prv->1))
                                                                 tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
     prv = --L.erase(prv);
                                                                 tp[1]=max(
   if (prv == L.begin() && now.a == prv->a)
                                                                  dp[u][0]+dp[v][0]+1
   L.erase(prv);
                                                                  dp[u][1]+max(dp[v][0],dp[v][1])
  if (it != L.end())
                                                                }else{
   while (it != --L.end() &&
                                                                 tp[0]=dp[u][0]+dp[v][0];
    (*it)(it->r) <= now(it->r))
                                                                 tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
     it = L.erase(it)
  if (it != L.begin()) {
                                                                dp[u][0]=tp[0],dp[u][1]=tp[1];
                                                               }
  prv = prev(it);
   const_cast<Line*>(&*prv)->r=now.l=((*prv)&now);
                                                              }else{
                                                               for(int i=0;i<(int)g[u].size();i++){</pre>
  if (it != L.end())
                                                                int v=g[u][i];
   const_cast<Line*>(&*it)->l=now.r=((*it)&now);
                                                                if(v==fa) continue;
  L.insert(it, now);
                                                                dfs(v,u);
11d Query(11d a) const { // query max at x=a
                                                               min_dp[0][0]=0;
                                                               min_dp[1][1]=1;
 if (L.empty()) return -INF;
                                                               min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
 Line::flag = false;
  auto it = --L.upper_bound({0, 0, a, 0});
                                                               for(int i=0;i<(int)g[u].size();i++){</pre>
  return (*it)(a);
                                                                int v=g[u][i];
                                                                if(v==fa) continue;
};
                                                                memset(tmp,0x8f,sizeof tmp);
                                                                tmp[0][0]=max(
8.6 Josephus Problem
                                                                 \min_{dp[0][0]+\max(dp[v][0],dp[v][1])}
// n people kill m for each turn
                                                                 min_dp[0][1]+dp[v][0]
int f(int n, int m) {
int s = 0;
                                                                tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
for (int i = 2; i <= n; i++)
                                                                tmp[1][0]=max(
 s = (s + m) \% i;
                                                                 \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
                                                                 min_dp[1][1]+dp[v][0]
return s;
```

```
ans[d] = row[i];
                                                                     for(int j = R[i]; j != i; j = R[j])
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
                                                                      remove(col[j]);
   memcpy(min_dp,tmp,sizeof tmp);
                                                                     dance(d+1);
 dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
                                                                     for(int j = L[i]; j != i; j = L[j])
  dp[u][0]=min_dp[0][0];
                                                                      resume(col[j]);
                                                                   resume(c);
int main(){
int m,a,b;
                                                                 } sol;
scanf("%d%d",&n,&m);
                                                                 8.9
                                                                       Tree Knapsack
for(int i=0;i<m;i++){</pre>
 scanf("%d%d",&a,&b);
                                                                 int dp[N][K];PII obj[N];
  init_g[a].push_back(b);
                                                                 vector<int> G[N];
 init_g[b].push_back(a);
                                                                 void dfs(int u,
                                                                  for(int s: G[u]) {
                                                                   if(mx < obj[s].first) continue;
for(int i=0;i<=mx-obj[s].FF;i++)</pre>
par[1]=-1;
tarjan(1);
dfs(1,-1);
                                                                     dp[s][i] = dp[u][i];
printf("%d\n", max(dp[1][0], dp[1][1]));
                                                                   dfs(s, mx - obj[s].first);
return 0;
                                                                   for(int i=obj[s].FF;i<=mx;i++)</pre>
                                                                     dp[u][i] = max(dp[u][i],
                                                                      dp[s][i - obj[s].FF] + obj[s].SS);
8.8 DLX
struct DLX {
const static int maxn=210;
                                                                 int main(){
 const static int maxm=210;
                                                                  int n, k; cin >> n >> k;
const static int maxnode=210*210;
                                                                  for(int i=1;i<=n;i++){</pre>
int n, m, size, row[maxnode], col[maxnode];
int U[maxnode], D[maxnode], L[maxnode], R[maxnode];
                                                                   int p; cin >> p;
                                                                   G[p].push_back(i);
int H[maxn], S[maxm], ansd, ans[maxn];
                                                                   cin >> obj[i].FF >> obj[i].SS;
 void init(int _n, int _m) {
 n = _n, m = _m;
                                                                  dfs(0, k); int ans = 0;
 for(int i = 0; i <= m; ++i) {</pre>
                                                                  for(int i=0;i<=k;i++) ans = max(ans, dp[0][i]);</pre>
   S[i] = 0;
                                                                  cout << ans << '\n';
   U[i] = D[i] = i;
                                                                  return 0;
   L[i] = i-1, R[i] = i+1;
 R[L[0] = size = m] = 0;
                                                                 8.10 N Queens Problem
  for(int i = 1; i <= n; ++i) H[i] = -1;
                                                                 vector< int > solve( int n ) {
                                                                  // no solution when n=2, 3
void Link(int r, int c) {
 ++S[col[++size] = c];
                                                                  vector< int > ret;
                                                                  if ( n % 6 == 2 ) {
  for ( int i = 2 ; i <= n ; i += 2 )
    ret.push_back( i );</pre>
  row[size] = r; D[size] = D[c];
 U[D[c]] = size; U[size] = c; D[c] = size;
  if(H[r] < 0) H[r] = L[size] = R[size] = size;
                                                                   ret.push_back( 3 ); ret.push_back( 1 );
 else {
                                                                   for ( int i = 7 ; i <= n ; i += 2 )
   R[size] = R[H[r]];
                                                                    ret.push_back( i );
   L[R[H[r]]] = size;
                                                                   ret.push_back( 5 );
   L[size] = H[r];
                                                                  } else if ( n % 6 == 3 ) {
for ( int i = 4 ; i <= n ; i += 2 )</pre>
   R[H[r]] = size;
                                                                    ret.push_back( i );
}
                                                                   ret.push_back( 2 );
void remove(int c) {
                                                                   for ( int i = 5 ; i <= n ; i += 2 )
 L[R[c]] = L[c]; R[L[c]] = R[c];
for(int i = D[c]; i != c; i = D[i])
                                                                    ret.push_back( i );
                                                                   ret.push_back( 1 ); ret.push_back( 3 );
   for(int j = R[i]; j != i; j = R[j]) {
                                                                  } else {
    U[D[j]] = U[j];
D[U[j]] = D[j];
                                                                   for ( int i = 2 ; i <= n ; i += 2 )
                                                                    ret.push_back( i );
    --S[col[j]];
                                                                   for ( int i = 1 ; i <= n ; i += 2 )
                                                                    ret.push_back( i );
void resume(int c) {
                                                                  return ret;
 L[R[c]] = c; R[L[c]] = c;
for(int i = U[c]; i != c; i = U[i])
                                                                 }
   for(int j = L[i]; j != i; j = L[j]) {
                                                                 8.11 Aliens Optimization
    U[D[j]] = j;
                                                                 long long Alien()
    D[U[j]] =
                                                                  long long c = kInf;
    ++S[col[j]];
                                                                  for (int d = 60; d >= 0; --d) {
   }
                                                                   // cost can be negative, depending on the problem.
}
                                                                   if (c - (1LL << d) < 0) continue;</pre>
void dance(int d) {
                                                                   long long ck = c - (1LL \ll d);
 if(d>=ansd) return;
                                                                   pair<long long, int> r = check(ck);
if (r.second == k) return r.first - ck * k;
 if(R[0] == 0) {
  ansd = d;
                                                                   if (r.second < k) c = ck;
   return;
                                                                  pair<long long, int> r = check(c);
 int c = R[0];
                                                                  return r.first - c * k;
  for(int i = R[0]; i; i = R[i])
  if(S[i] < S[c]) c = i;
  remove(c);
  for(int i = D[c]; i != c; i = D[i]) {
```