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```
is nu bs=2 ru mouse=a encoding=utf-8 ls=2
cin cino+=j1 et sw=4 sts=4 tgc sc hls
n on
lorscheme desert
letype indent on
oremap {<CR> {<CR>}<ESC>0
o <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -std=c++17 -
 DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
 Wconversion -fsanitize=address,undefined -g && echo
  success<CR>
o <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 &&
  echo success<CR>
o <F10> <ESC>:!./"%<"<CR>
```

Debug Macro

```
def KISEKI
efine safe cerr<<__PRETTY_FUNCTION__\
<<" line "<<__LINE__<<" safe\n"
efine debug(a...) qwerty(#a, a)
efine orange(a...) dvorak(#a, a)
ing std::cerr;
nplate <typename ...T>
id qwerty(const char *s, T ...a) {
cerr << "\e[1;32m(" << s << ") = (";
int cnt = sizeof...(T);
..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
nplate <typename Iter>
id dvorak(const char *s, Iter L, Iter R) {
cerr << "\e[1;32m[ " << s << " ] = [ ";
efine safe ((void)0)
efine debug(...) ((void)0)
efine orange(...) ((void)0)
ndif
```

Increase Stack

```
nst int size = 256 << 20;
                                       gister long rsp asm("rsp<sup>"</sup>);
                                       r *p = (char*)malloc(size)+size, *bak = (char*)rsp;
                                       asm__("mòvq %0́, %%rsp\n"::"r"(p));
                                       main
```

Pragma Optimization

```
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
1.5 IO Optimization
static inline int gc() {
 constexpr int B = 1<<20;</pre>
 static char buf[B], *p, *q;
 if(p == q \&\&
  (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
  return EOF:
 return *p++;
template < typename T >
static inline bool gn( T &x ) {
 int c = gc(); T sgn = 1; x = 0;
while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
if(c == '-') sgn = -1, c = gc();
 if(c == EOF) return false;
 while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
 return x *= sgn, true;
```

#pragma GCC optimize("Ofast,no-stack-protector")

#pragma GCC optimize("no-math-errno,unroll-loops")

Data Structure 2

Dark Magic 2.1

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using namespace __gnu_pbds;
// heap tags: paring/binary/binomial/rc_binomial/thin
template<typename T>
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>, \
                  pairing_heap_tag>;
// pbds_heap::point_iterator
// x = pq.push(10); pq.modify(x, 87); a.join(b);
// tree tags: rb_tree_tag/ov_tree_tag/splay_tree_tag
template<typename T>
using ordered_set = tree<T, null_type, less<T>
   rb_tree_tag, tree_order_statistics_node_update>;
// find_by_order, order_of_key
// hash tables: cc_hash_table/gp_hash_table
```

2.2 Link-Cut Tree

p->ch[dir]=c;

```
struct Node{
Node *par, *ch[2];
int xor_sum, v;
bool is_rev;
Node(int _v){
 v=xor_sum=_v;is_rev=false;
 par=ch[0]=ch[1]=nullptr;
inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
inline void down(){
 if(is_rev){
  if(ch[0]!=nullptr) ch[0]->set_rev();
   if(ch[1]!=nullptr) ch[1]->set_rev();
   is_rev=false;
 }
inline void up(){
 xor_sum=v;
  if(ch[0]!=nullptr){
  xor_sum^=ch[0]->xor_sum;
  ch[0]->par=this;
 if(ch[1]!=nullptr){
  xor_sum^=ch[1]->xor_sum;
  ch[1]->par=this;
inline bool is_root(){
 return par==nullptr ||\
   (par->ch[0]!=this && par->ch[1]!=this);
bool is_rch(){return !is_root() && par->ch[1]==this;}
} *node[maxn], *stk[maxn];
int top;
void to_child(Node* p,Node* c,bool dir){
```

```
p->up();
inline void rotate(Node* node){
 Node* par=node->par;
 Node* par_par=par->par;
 bool dir=node->is_rch()
 bool par_dir=par->is_rch()
 to_child(par, node->ch[!dir], dir);
 to_child(node,par,!dir);
 if(par_par!=nullptr && par_par->ch[par_dir]==par)
  to_child(par_par,node,par_dir);
 else node->par=par_par;
inline void splay(Node* node){
 Node* tmp=node;
 stk[top++]=node;
 while(!tmp->is_root()){
  tmp=tmp->par;
  stk[top++]=tmp;
 while(top) stk[--top]->down();
 for(Node *fa=node->par;
  !node->is_root();
  rotate(node), fa=node->par)
  if(!fa->is_root())
   rotate(fa->is_rch()==node->is_rch()?fa:node);
inline void access(Node* node){
 Node* last=nullptr;
 while(node!=nullptr){
  splay(node);
  to_child(node, last, true);
  last=node;
  node=node->par;
inline void change_root(Node* node){
 access(node);splay(node);node->set_rev();
inline void link(Node* x, Node* y){
 change_root(x);splay(x);x->par=y;
inline void split(Node* x,Node* y){
 change_root(x);access(y);splay(x);
 to_child(x,nullptr,true);y->par=nullptr;
inline void change_val(Node* node,int v){
access(node);splay(node);node->v=v;node->up();
inline int query(Node* x,Node* y){
 change_root(x);access(y);splay(y);
 return y->xor_sum;
inline Node* find_root(Node* node){
 access(node);splay(node);
 Node* last=nullptr:
 while(node!=nullptr){
  node->down();last=node;node=node->ch[0];
 return last;
set<pii> dic;
inline void add_edge(int u,int v){
 if(u>v) swap(u,v)
 if(find_root(node[u])==find_root(node[v])) return;
 dic.insert(pii(u,v))
link(node[u],node[v]);
inline void del_edge(int u,int v){
 if(u>v) swap(u,v);
 if(dic.find(pii(u,v))==dic.end()) return;
 dic.erase(pii(u,v))
 split(node[u],node[v]);
2.3 LiChao Segment Tree
 int m, k, id;
 Line() : id( -1 ) {}
```

```
struct Line{
 Line('int a, int'b,'int c')
: m(a), k(b), id(c) {}
 int at( int x ) { return m * x + k; }
```

```
#undef sz
class LiChao {
 private:
                                                             2.5 Linear Basis
  int n; vector< Line > nodes;
  inline int lc( int x ) { return 2 * x + 1; }
                                                             template <int BITS>
  inline int rc( int x ) { return 2 * x + 2; }
                                                             struct LinearBasis {
  void insert( int 1, int r, int id, Line ln ) {
                                                              array<uint64_t, BITS> basis;
   int m = (1 + r) >> 1;
                                                              Basis() { basis.fill(0); }
                                                              void add(uint64_t x)
   if ( nodes[ id ].id == -1 ) {
    nodes[ id ] = ln;
                                                               for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
                                                                if (basis[i] == 0) {
    return:
                                                                 basis[i] = x;
   bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
                                                                 return;
   if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
    atLeft ^= 1; swap( nodes[ id ], ln );
                                                                x ^= basis[i];
                                                               }
   if ( r - 1 == 1 ) return;
   if ( atLeft ) insert( l, m, lc( id ), ln );
                                                              bool ok(uint64_t x) {
   else insert( m, r, rc( id ), ln );
                                                               for (int i = 0; i < BITS; ++i)</pre>
                                                                if ((x >> i) & 1) x ^= basis[i];
  int query( int 1, int r, int id, int x ) {
                                                               return x == 0;
   int ret = 0;
   if ( nodes[ id ].id != -1 )
                                                             };
    ret = nodes[ id ].at( x );
                                                             2.6
                                                                   Binary Search On Segment Tree
   int m = (1 + r) >> 1;
   if ( r - l == 1 ) return ret;
                                                             // find_first = x -> minimal x s.t. check( [a, x) )
   else if (x < m )
                                                             // find_last = x \rightarrow maximal x s.t. check([x, b))
    return max( ret, query( 1, m, lc( id ), x ) );
                                                             template <typename C>
   else
                                                             int find_first(int 1, const C &check) {
    return max( ret, query( m, r, rc( id ), x ) );
                                                              if (1 >= n) return n;
                                                              1 += sz;
 public:
                                                              for (int i = height; i > 0; i--)
  void build( int n_ ) {
                                                               propagate(l >> i);
  n = n_; nodes.clear();
                                                              Monoid sum = identity;
   nodes.resize( n << 2, Line() );</pre>
                                                               while ((1 & 1) == 0) 1 >>= 1;
  void insert( Line ln ) { insert( 0, n, 0, ln ); }
                                                               if (check(f(sum, data[1]))) {
  int query( int x ) { return query( 0, n, 0, x ); }
                                                                while (1 < sz) {</pre>
                                                                 propagate(1);
                                                                 1 <<= 1;
2.4 Treap
                                                                 auto nxt = f(sum, data[1]);
namespace Treap{
                                                                 if (not check(nxt)) {
 #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                                  sum = nxt;
 struct node{
                                                                  1++;
  int size;
                                                                 }
  uint32_t pri;
                                                                }
  node *lc, *rc, *pa;
                                                                return 1 + 1 - sz;
  node():size(0),pri(rand()),lc(0),rc(0),pa(0){}
  void pull() {
                                                               sum = f(sum, data[1++]);
  size = 1; pa = nullptr;
                                                              } while ((1 & -1) != 1);
   if ( lc ) { size += lc->size; lc->pa = this; }
if ( rc ) { size += rc->size; rc->pa = this; }
                                                              return n;
  }
                                                             template <typename C>
                                                             int find_last(int r, const C &check) {
node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
                                                              if (r <= 0) return -1;
                                                              r += sz;
  if ( L->pri > R->pri ) {
                                                              for (int i = height; i > 0; i--)
  L->rc = merge( L->rc, R ); L->pull();
                                                               propagate((r - 1) >> i);
   return L;
                                                              Monoid sum = identity;
  } else {
                                                              do {
   R->lc = merge( L, R->lc ); R->pull();
   return R;
                                                               while (r > 1 \text{ and } (r \& 1)) r >>= 1;
  }
                                                               if (check(f(data[r], sum))) {
                                                                while (r < sz) {</pre>
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                                 propagate(r);
  if ( not rt ) L = R = nullptr;
                                                                 r = (r << 1) + 1;
  else if( sz( rt->lc ) + 1 <= k ) {
                                                                 auto nxt = f(data[r], sum);
                                                                 if (not check(nxt)) {
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                  sum = nxt;
   L->pull();
                                                                  r--:
  } else {
                                                                 }
   R = rt
   split_by_size( rt->lc, k, L, R->lc );
                                                                return r - sz;
   R->pull();
                                                               sum = f(data[r], sum);
  }
                                                              } while ((r & -r) != r);
 int getRank(node *o) {
                                                              return -1;
  int r = sz(o->lc);
  for (;o->pa != nullptr; o = o->pa)
   if (o->pa->rc != o) r += sz(o->pa->lc);
                                                                  Graph
  return r;
```

3.1 2-SAT (SCC)

G.clear(); G.resize(n = n_);

```
class TwoSat{
                                                                 low.assign(n, ecnt = 0);
private:
                                                                dfn.assign(n, 0);
 int n;
                                                               }
  vector<vector<int>> rG,G,sccs;
                                                               void add_edge(int u, int v) {
                                                                G[u].emplace_back(v, ecnt);
 vector<int> ord,idx;
  vector<bool> vis,result;
                                                                G[v].emplace_back(u, ecnt++);
  void dfs(int u){
                                                               void solve() {
   vis[u]=true
   for(int v:G[u])
                                                                bridge.assign(ecnt, false);
                                                                for (int i = 0; i < n; ++i)
    if(!vis[v]) dfs(v);
                                                                 if (not dfn[i]) dfs(i, i);
   ord.push_back(u);
  void rdfs(int u){
                                                               bool is_bridge(int x) { return bridge[x]; }
  vis[u]=false;idx[u]=sccs.size()-1;
                                                             } bcc_bridge;
   sccs.back().push_back(u);
                                                             3.3 BCC Vertex
   for(int v:rG[u])
    if(vis[v])rdfs(v);
                                                             class BCC_AP {
                                                              private:
public:
                                                               int n, ecnt;
  void init(int n_){
                                                               vector<vector<pair<int,int>>> G;
  n=n_;G.clear();G.resize(n);
                                                               vector<int> bcc, dfn, low, st;
                                                               vector<bool> ap, ins;
void dfs(int u, int f)
   rG.clear();rG.resize(n);
   sccs.clear();ord.clear();
                                                                dfn[u] = low[u] = dfn[f] + 1;
   idx.resize(n);result.resize(n);
                                                                 int ch = 0;
                                                                for (auto [v, t]: G[u]) if (v != f) {
  if (not ins[t]) {
  void add_edge(int u,int v){
  G[u].push_back(v);rG[v].push_back(u);
                                                                  st.push_back(t);
                                                                   ins[t] = true;
  void orr(int x,int y){
   if ((x^y)==1)return
   add_edge(x^1,y); add_edge(y^1,x);
                                                                  if (dfn[v]) {
                                                                  low[u] = min(low[u], dfn[v]);
  bool solve(){
                                                                   continue:
   vis.clear();vis.resize(n);
                                                                  } ++ch; dfs(v, u);
   for(int i=0;i<n;++i)</pre>
                                                                  low[u] = min(low[u], low[v]);
    if(not vis[i])dfs(i);
                                                                  if (low[v] >= dfn[u]) {
                                                                  ap[u] = true;
   reverse(ord.begin(),ord.end());
   for (int u:ord){
                                                                   while (true) {
    if(!vis[u])continue;
                                                                    int eid = st.back(); st.pop_back();
                                                                    bcc[eid] = ecnt;
    sccs.push_back(vector<int>());
    rdfs(u);
                                                                    if (eid == t) break;
   for(int i=0;i<n;i+=2)</pre>
                                                                  ecnt++;
    if(idx[i]==idx[i+1])
                                                                 }
     return false:
   vector<bool> c(sccs.size());
                                                                 if (ch == 1 and u == f) ap[u] = false;
   for(size_t i=0;i<sccs.size();++i){</pre>
    for(size_t j=0;j<sccs[i].size();++j){</pre>
                                                              public:
     result[sccs[i][j]]=c[i];
                                                               void init(int n_) {
     c[idx[sccs[i][j]^1]]=!c[i];
                                                                G.clear(); G.resize(n = n_);
   }
                                                                 ecnt = 0; ap.assign(n, false);
                                                                low.assign(n, 0); dfn.assign(n, 0);
   return true;
                                                               void add_edge(int u, int v) {
                                                                G[u].emplace_back(v, ecnt);
G[v].emplace_back(u, ecnt++);
  bool get(int x){return result[x];}
  inline int get_id(int x){return idx[x];}
  inline int count(){return sccs.size();}
                                                               void solve() {
                                                                ins.assign(ecnt, false);
3.2 BCC Edge
                                                                bcc.resize(ecnt); ecnt = 0;
                                                                for (int i = 0; i < n; ++i)
if (not dfn[i]) dfs(i, i);</pre>
class BCC_Bridge {
private:
 int n, ecnt;
                                                               int get_id(int x) { return bcc[x]; }
  vector<vector<pair<int,int>>> G;
  vector<int> dfn, low;
                                                               int count() { return ecnt;
  vector<bool> bridge;
                                                               bool is_ap(int x) { return ap[x]; }
  void dfs(int u, int f)
                                                             } bcc_ap;
   dfn[u] = low[u] = dfn[f] + 1;
                                                             3.4 Centroid Decomposition
   for (auto [v, t]: G[u]) {
    if (v == f) continue;
                                                             struct Centroid {
    if (dfn[v]) {
                                                              vector<vector<int64_t>> Dist;
                                                              vector<int> Parent, Depth;
     low[u] = min(low[u], dfn[v]);
     continue;
                                                              vector<int64_t> Sub, Sub2;
                                                              vector<int> Sz, Sz2;
    dfs(v, u);
                                                              Centroid(vector<vector<pair<int, int>>> g) {
    low[u] = min(low[u], low[v]);
                                                               int N = g.size();
    if (low[v] > dfn[u]) bridge[t] = true;
                                                               vector<bool> Vis(N);
                                                               vector<int> sz(N), mx(N);
                                                               vector<int> Path;
                                                               Dist.resize(N)
public:
  void init(int n_) {
                                                               Parent.resize(N);
```

Depth.resize(N);

in[e.v] = e.w;

```
auto DfsSz = [&](auto dfs, int x) -> void {
                                                                   prv[e.v] = e.u;
   Vis[x] = true; sz[x] = 1; mx[x] = 0;
   for (auto [u, w] : g[x]) {
                                                                 in[root] = 0;
    if (Vis[u]) continue;
                                                                 prv[root] = -1;
    dfs(dfs, u)
                                                                 for (int i = 0; i < n; i++)
                                                                  if (in[i] == -inf)
    sz[x] += sz[u];
    mx[x] = max(mx[x], sz[u]);
                                                                   return -inf;
                                                                  // find cycle
   Path.push_back(x);
                                                                  int tot = 0;
                                                                 vector<int> id(n, -1), vis(n, -1);
for (int i = 0; i < n; i++) {</pre>
  }:
  auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
                                                                  ans += in[i];
   Dist[x].push_back(D);Vis[x] = true;
                                                                  for (int x = i; x != -1 && id[x] == -1; x = prv[x])
   for (auto [u, w] : g[x]) {
   if (Vis[u]) continue;
                                                                   if (vis[x] == i) {
    dfs(dfs, u, D + w);
                                                                     for (int y = prv[x]; y != x; y = prv[y])
                                                                      id[y] = tot;
  };
                                                                     id[x] = tot++;
  auto Dfs = [&]
                                                                    break;
   (auto dfs, int x, int D = 0, int p = -1)->void {
   Path.clear(); DfsSz(DfsSz, x);
                                                                   vis[x] = i;
   int M = Path.size();
                                                                  }
   int C = -1;
   for (int u : Path) {
                                                                 if (!tot)
                                                                  return ans;
    if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
    Vis[u] = false;
                                                                 for (int i = 0; i < n; i++)</pre>
                                                                  if (id[i] == -1)
   DfsDist(DfsDist, C);
                                                                   id[i] = tot++;
   for (int u : Path) Vis[u] = false;
                                                                  // shrink
   Parent[C] = p; Vis[C] = true;
                                                                 for (auto &e : E) {
   Depth[C] = D;
                                                                  if (id[e.u] != id[e.v])
   for (auto [u, w] : g[C]) {
                                                                   e.w -= in[e.v];
    if (Vis[u]) continue
                                                                  e.u = id[e.u], e.v = id[e.v];
    dfs(dfs, u, D + 1, C);
                                                                 n = tot;
                                                                 root = id[root];
  Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
  Sz.resize(N); Sz2.resize(N);
                                                                assert(false);
void Mark(int v) {
                                                              } DMST:
  int x = v, z = -1
                                                              3.6 Dominator Tree
 for (int i = Depth[v]; i >= 0; --i) {
Sub[x] += Dist[v][i]; Sz[x]++;
                                                              namespace dominator {
  if (z != -1) {
                                                              vector<int> g[maxn], r[maxn], rdom[maxn];
                                                              int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
    Sub2[z] += Dist[v][i];
    Sz2[z]++;
                                                              void init(int n) {
   z = x; x = Parent[x];
                                                               // vertices are numbered from 0 to n-1
  }
                                                               fill(dfn, dfn + n, -1);fill(rev, rev + n, -1);
                                                               fill(fa, fa + n, -1); fill(val, val + n, -1);
int64_t Query(int v) {
                                                               fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
                                                               fill(dom, dom + n, -1); tk = 0;
for (int i = 0; i < n; ++i) {
 int64_t res = 0;
 int x = v, z = -1;
 for (int i = Depth[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
                                                                g[i].clear(); r[i].clear(); rdom[i].clear();
  if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
  z = x; x = Parent[x];
                                                              void add_edge(int x, int y) { g[x].push_back(y); }
                                                              void dfs(int x) {
                                                               rev[dfn[x] = tk] = x;
  return res;
                                                               fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
                                                               for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
};
3.5 Directed Minimum Spanning Tree
                                                                r[dfn[u]].push_back(dfn[x]);
struct DirectedMST { // find maximum
                                                               }
struct Edge {
                                                              void merge(int x, int y) { fa[x] = y; }
 int u, v;
                                                              int find(int x, int c = 0) {
  int w;
                                                               if (fa[x] == x) return c ? -1 : x;
 Edge(int u, int v, int w) : u(u), v(v), w(w) {}
                                                               int p = find(fa[x], 1);
                                                               if (p == -1) return c ? fa[x] : val[x];
vector<Edge> Edges;
                                                               if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
void clear() { Edges.clear(); }
void addEdge(int a, int b, int w) { Edges.emplace_back
                                                               fa[x] = p
                                                               return c ? p : val[x];
    (a, b, w); }
int solve(int root, int n) {
  vector<Edge> E = Edges;
                                                              vector<int> build(int s, int n) {
                                                              // return the father of each node in the dominator tree
  int ans = 0:
 while (true) {
                                                              // p[i] = -2 if i is unreachable from s
   // find best in edge
                                                               dfs(s);
   vector<int> in(n, -inf), prv(n, -1);
                                                               for (int i = tk - 1; i >= 0; --i) {
                                                                for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
                                                                if (i) rdom[sdom[i]].push_back(i);
```

for (int &u : rdom[i]) {

```
if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
   int p = find(u);
   if (sdom[p] == i) dom[u] = i;
                                                                       chain[ u ] = chain[ v ];
   else dom[u] = p;
                                                                     if ( not chain[ u ] )
  if (i) merge(i, rp[i]);
                                                                      chain[ u ] = chain_ ++;
vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)</pre>
                                                                    void dfschain( int u, int f ) {
                                                                     fa[ u ][ 0 ] = f;
for ( int i = 1 ; i < LOG_N ; ++ i )</pre>
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
                                                                     fa[u][i] = fa[fa[u][i-1]][i-1];
tl[u] = time_++;
 for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
 return p;
                                                                     if ( not chain_st[ chain[ u ] ] )
                                                                      chain_st[ chain[ u ] ] = u;
3.7 Edge Coloring
                                                                     for ( int v : G[ u ]
                                                                      if ( v != f and chain[ v ] == chain[ u ] )
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
                                                                       dfschain( v, u );
                                                                     for ( int v : G[ u ] )
void clear(int N) {
                                                                      if ( v != f and chain[ v ] != chain[ u ] )
 for (int i = 0; i <= N; i++)
  for (int j = 0; j <= N; j++)
C[i][j] = G[i][j] = 0;
                                                                       dfschain( v, u );
                                                                     tr[ u ] = time_;
                                                                    bool anc( int u, int v ) {
  return tl[ u ] <= tl[ v ] and tr[ v ] <= tr[ u ];</pre>
void solve(vector<pair<int, int>> &E, int N) {
int X[kN] = {}, a;
auto update = [&](int u) {
  for (X[u] = 1; C[u][X[u]]; X[u]++);
                                                                   public:
                                                                    int lca( int u, int v ) {
  if ( anc( u, v ) ) return u;
 auto color = [&](int u, int v, int c) {
  int p = G[u][v];
                                                                     for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
                                                                      if ( not anc( fa[ u ][ i ], v ) )
u = fa[ u ][ i ];
  G[u][v] = G[v][u] = c;
  C[u][c] = v, C[v][c] = u;
  C[u][p] = C[v][p] = 0;
                                                                     return fa[ u ][ 0 ];
  if (p) X[u] = X[v] = p;
                                                                    void init( int n ) {
  else update(u), update(v);
                                                                     fa.assign( ++n, vector< int >( LOG_N ) );
  return p;
                                                                     for ( LOG_N = 0; ( 1 << LOG_N) < n; ++ LOG_N);
 };
                                                                     G.clear(); G.resize( n );
tl.assign( n, 0 ); tr.assign( n, 0 );
 auto flip = [&](int u, int c1, int c2) {
  int p = C[u][c1];
  swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
if (!C[u][c1]) X[u] = c1;
                                                                     chain.assig( n, 0 ); chain_st.assign( n, 0 );
                                                                    void add_edge( int u , int v ) {
  if (!C[u][c2]) X[u] = c2;
                                                                     // 1-base
                                                                     G[ u ].push_back( v );
  return p;
                                                                     G[ v ].push_back( u );
for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {
  auto [u, v] = E[t];</pre>
                                                                    void decompose(){
                                                                     chain_ = 1;
  int v0 = v, c = X[u], c0 = c, d;
                                                                     predfs( 1, 1 );
                                                                     time_{-} = 0;
  vector<pair<int, int>> L; int vst[kN] = {};
  while (!G[u][v0]) {
                                                                     dfschain(1,1);
   L.emplace_back(v, d = X[v]);
   if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
                                                                    PII get_subtree(int u) { return {tl[ u ],tr[ u ] }; }
     c = color(u, L[a].first, c);
                                                                    vector< PII > get_path( int u , int v ){
                                                                     vector< PII > res;
   else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
                                                                     int g = lca( u, v );
while ( chain[ u ] != chain[ g ] ) {
  int s = chain_st[ chain[ u ] ];
     color(u, L[a].first, L[a].second);
   else if (vst[d]) break
   else vst[d] = 1, v = C[u][d];
                                                                      res.emplace_back( tl[ s ], tl[ u ] + 1 );
  if (!G[u][v0]) {
  for (; v; v = flip(v, c, d), swap(c, d));
                                                                      u = fa[ s ][ 0 ];
                                                                     res.emplace_back( tl[ g ], tl[ u ] + 1 );
while ( chain[ v ] != chain[ g ] ) {
  int s = chain_st[ chain[ v ] ];
   if (C[u][c0]) { a = int(L.size()) - 1;
    while (--a >= 0 && L[a].second != c);
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
                                                                      res.emplace_back( tl[ s ], tl[ v ] + 1 );
   } else t--;
                                                                      v = fa[ s ][ 0 ];
                                                                     res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
                                                                     return res;
     Lowbit Decomposition
                                                                     /* res : list of intervals from u to v
                                                                      \star ( note only nodes work, not edge )
class LowbitDecomp{
                                                                      * usage :
private:
                                                                      * vector< PII >& path = tree.get_path( u , v )
 int time_, chain_, LOG_N;
                                                                      * for( auto [ 1, r ] : path ) {
 vector< vector< int > > G, fa;
                                                                      * 0-base [ 1, r )
 vector< int > tl, tr, chain, chain_st;
 // chain_- : number of chain
                                                                      * }
                                                                      */
 // tl, tr[ u ] : subtree interval in the seq. of u
 // chain_st[ u ] : head of the chain contains u
 // chian[u] : chain id of the chain u is on
                                                                  } tree;
 void predfs( int u, int f ) {
                                                                   3.9 Manhattan Minimum Spanning Tree
  chain[ u ] = 0;
  for ( int v : G[ u ] ) {
  if ( v == f ) continue;
                                                                  typedef Point<int> P;
                                                                   vector<array<int, 3>> manhattanMST(vector<P> ps) {
   predfs( v, u );
                                                                   vi id(sz(ps));
```

```
iota(all(id), 0);
                                                                 for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                  deg[ i ] = G[ i ].count();
 vector<array<int, 3>> edges;
                                                                 bits pob, nob = 0; pob.set();
 rep(k, 0, 4) {
                                                                 for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
  sort(all(id), [&](int i, int j) {
                                                                 for ( size_t i = 0 ; i < n ; ++ i ) {
   return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
                                                                  size_t v = deo[ i ];
  map<int, int> sweep;
                                                                  bits tmp; tmp[ v ] = 1;
                                                                  BK( tmp, pob & G[ v ], nob & G[ v ] );
  for (int i : id) {
                                                                  pob[ v ] = 0, nob[ v ] = 1;
   for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++)) {
    int j = it->second
                                                                 return static_cast< int >( ans.count() );
    P d = ps[i] - ps[j];
    if (d.y > d.x) break;
                                                              };
    edges.push_back({d.y + d.x, i, j});
                                                               3.11 MaxCliqueDyn
                                                               constexpr int kN = 150;
   sweep[-ps[i].y] = i;
                                                               struct MaxClique { // Maximum Clique
                                                                bitset<kN> a[kN], cs[kN];
  for (P &p : ps)
   if (k \& 1) p.x = -p.x;
                                                                int ans, sol[kN], q, cur[kN], d[kN], n;
   else swap(p.x, p.y);
                                                                void init(int _n) {
                                                                 n = _n, ans = q = 0;
 return edges; // [{w, i, j}, ...]
                                                                 for (int i = 0; i < n; i++) a[i].reset();
                                                                void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
3.10
      MaxClique
                                                                void csort(vector<int> &r, vector<int> &c) {
                                                                 int mx = 1, km = max(ans - q + 1, 1), t = 0,
// contain a self loop u to u, than u won't in clique
template < size_t MAXN >
                                                                   m = int(r.size())
                                                                 cs[1].reset(); cs[2].reset();
class MaxClique{
                                                                 for (int i = 0; i < m; i++) {
private:
                                                                  int p = r[i], k = 1;
 using bits = bitset< MAXN >;
                                                                  while ((cs[k] & a[p]).count()) k++;
 bits popped, G[ MAXN ], ans;
 size_t deg[ MAXN ], deo[ MAXN ], n;
                                                                  if (k > mx) cs[++mx + 1].reset();
 void sort_by_degree() {
                                                                  cs[k][p] = 1;
                                                                  if (k < km) r[t++] = p;
  popped.reset();
  for ( size_t i = 0 ; i < n ; ++ i )
                                                                 c.resize(m);
if (t) c[t - 1] = 0;
    deg[ i ] = G[ i ].count();
  for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;
                                                                 for (int k = km; k <= mx; k++) {</pre>
    for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )</pre>
                                                                  for (int p = int(cs[k]._Find_first());
  p < kN; p = int(cs[k]._Find_next(p))) {</pre>
        mi = deg[ id = j ];
                                                                   r[t] = p; c[t++] = k;
    popped[ deo[ i ] = id ] = 1;
    for( size_t u = G[ i ]._Find_first() ;
                                                                 }
     u < n ; u = G[ i ]._Find_next( u ) )
                                                                void dfs(vector<int> &r, vector<int> &c, int 1,
      -- deg[ u ];
  }
                                                                 bitset<kN> mask) {
                                                                 while (!r.empty()) {
 void BK( bits R, bits P, bits X ) {
                                                                  int p = r.back(); r.pop_back();
  if (R.count()+P.count() <= ans.count()) return;</pre>
                                                                  mask[p] = 0;
  if ( not P.count() and not X.count() ) {
                                                                  if (q + c.back() <= ans) return;</pre>
                                                                  cur[q++] = p;
   if ( R.count() > ans.count() ) ans = R;
   return;
                                                                  vector<int> nr, nc;
                                                                  bitset<kN> nmask = mask & a[p];
  /* greedily chosse max degree as pivot
                                                                  for (int i : r)
  bits cur = P | X; size_t pivot = 0, sz = 0;
                                                                   if (a[p][i]) nr.push_back(i);
  for ( size_t u = cur._Find_first()
                                                                  if (!nr.empty()) {
  u < n ; u = cur._Find_next( u ) )</pre>
                                                                   if (1 < 4) {
   if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
                                                                    for (int i : nr)
  cur = P & ( ~G[ pivot ] );
                                                                     d[i] = int((a[i] & nmask).count());
  */ // or simply choose first
                                                                    sort(nr.begin(), nr.end(),
  bits cur = P & (~G[ ( P | X )._Find_first() ]);
                                                                     [&](int x, int y) {
  return d[x] > d[y];
  for ( size_t u = cur._Find_first()
   u < n ; u = cur._Find_next( u ) ) {</pre>
                                                                     });
   if ( R[ u ] ) continue;
                                                                  csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
   R[u] = 1;
   BK( R, P & G[ u ], X & G[ u ] );
   R[u] = P[u] = 0, X[u] = 1;
                                                                   ans = q; copy(cur, cur + q, sol);
                                                                  c.pop_back(); q--;
public:
                                                                 }
 void init( size_t n_ ) {
                                                                int solve(bitset<kN> mask) { // vertex mask
  n = n_{\perp}
                                                                 vector<int> r, c;
  for ( size_t i = 0 ; i < n ; ++ i )
                                                                 for (int i = 0; i < n; i++)
  if (mask[i]) r.push_back(i);</pre>
   G[ i ].reset();
  ans.reset();
                                                                 for (int i = 0; i < n; i++)</pre>
 void add_edges( int u, bits S ) { G[ u ] = S; }
void add_edge( int u, int v ) {
                                                                  d[i] = int((a[i] & mask).count());
                                                                 sort(r.begin(), r.end(),
 G[u][v] = G[v][u] = 1;
                                                                  [&](int i, int j) { return d[i] > d[j]; });
                                                                 csort(r, c);
 int solve() {
                                                                 dfs(r, c, 1, mask);
                                                                 return ans; // sol[0 ~ ans-1]
  sort_by_degree(); // or simply iota( deo... )
```

```
}
                                                                     dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
} graph;
                                                                    void shortest_path(){
3.12 Minimum Mean Cycle
                                                                     for( int k = 0 ; k < n ; k ++ )
for( int i = 0 ; i < n ; i ++ )
/* minimum mean cycle O(VE) */
                                                                       for( int j = 0 ; j < n ; j ++ )
dst[ i ][ j ] = min( dst[ i ][ j ],</pre>
struct MMC{
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
                                                                            dst[ i ][ k ] + dst[ k ][ j ] );
#define V 1021
#define inf 1e9
                                                                    int solve( const vector<int>& ter ){
 struct Edge { int v,u; double c; };
                                                                     int t = (int)ter.size();
 int n, m, prv[V][V], prve[V][V], vst[V];
                                                                     for( int i = 0 ; i < ( 1 << t ) ; i ++ )</pre>
                                                                      for( int j = 0 ; j < n ; j ++ )
dp[ i ][ j ] = INF;
 Edge e[E];
 vector<int> edgeID, cycle, rho;
 double d[V][V];
                                                                     for( int i = 0 ; i < n ; i ++ )</pre>
 void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
                                                                      dp[0][i] = 0;
                                                                     for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
                                                                      if( msk == ( msk & (-msk) ) ){
 void add_edge( int vi , int ui , double ci )
 { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
                                                                       int who = __lg( msk );
                                                                       for( int i = 0 ; i < n ; i ++ )
dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];</pre>
  for(int i=0; i<n; i++) d[0][i]=0;</pre>
  for(int i=0; i<n; i++) {</pre>
                                                                       continue:
   fill(d[i+1], d[i+1]+n, inf);
for(int j=0; j<m; j++) {
                                                                      for( int i = 0 ; i < n ; i ++ )</pre>
    int v = e[j].v, u = e[j].u;
                                                                       for( int submsk = ( msk - 1 ) & msk ; submsk ;
                                                                          submsk = ( submsk - 1 ) & msk )
dp[ msk ][ i ] = min( dp[ msk ][ i ],
    if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
     d[i+1][u] = d[i][v]+e[j].c;
     prv[i+1][u] = v;
                                                                                   dp[ submsk ][ i ] +
                                                                                   dp[ msk ^ submsk ][ i ] );
     prve[i+1][u] = j;
                                                                      for( int i = 0 ; i < n ; i ++ ){</pre>
                                                                       tdst[ i ] = INF;
                                                                       }
 double solve(){
  // returns inf if no cycle, mmc otherwise
                                                                      for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = tdst[ i ];</pre>
  double mmc=inf;
  int st = -1;
  bellman_ford();
  for(int i=0; i<n; i++) {
  double avg=-inf;</pre>
                                                                     int ans = INF;
                                                                     for( int i = 0 ; i < n ; i ++ )</pre>
   for(int k=0; k<n; k++) {</pre>
                                                                      ans = min( ans , dp[ ( 1 << t ) - 1 ][ i ] );
    if(d[n][i]<inf-eps)</pre>
                                                                     return ans;
     avg=max(avg,(d[n][i]-d[k][i])/(n-k));
                                                                    }
    else avg=max(avg,inf);
                                                                  } solver;
                                                                   3.14 Mo's Algorithm on Tree
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
                                                                   int q; vector< int > G[N];
  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
                                                                   struct Que{
  for (int i=n; !vst[st]; st=prv[i--][st]) {
                                                                    int u, v, id;
                                                                   } que[ N ];
   vst[st]++;
                                                                   int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
   edgeID.PB(prve[i][st]);
                                                                   void dfs( int u, int f ) {
  dfn[ u ] = dfn_++; int saved_rbp = stk_;
   rho.PB(st);
  while (vst[st] != 2) {
                                                                    for ( int v : G[ u ] ) {
                                                                     if ( v == f ) continue;
   int v = rho.back(); rho.pop_back();
   cycle.PB(v);
                                                                     dfs( v, u );
   vst[v]++;
                                                                     if ( stk_ - saved_rbp < SQRT_N ) continue;</pre>
                                                                     for ( ++ block_ ; stk_ != saved_rbp ; )
  block_id[ stk[ -- stk_ ] ] = block_;
  reverse(ALL(edgeID));
  edgeID.resize(SZ(cycle));
                                                                    stk[ stk_+ + ] = u;
  return mmc;
 }
} mmc;
                                                                   bool inPath[ N ];
                                                                   void Diff( int u ) {
  if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
3.13 Minimum Steiner Tree
// Minimum Steiner Tree
                                                                    else { /*add this edge*/ }
// 0(V 3^T + V^2 2^T)
struct SteinerTree{
                                                                   void traverse( int& origin_u, int u ) {
                                                                    for ( int g = lca( origin_u, u )
#define V 33
#define T 8
                                                                     origin_u != g ; origin_u = parent_of[ origin_u ] )
#define INF 1023456789
                                                                      Diff( origin_u );
 int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
                                                                    for (int v = u; v != origin_u; v = parent_of[v])
 void init( int _n ){
                                                                     Diff( v );
 n = _n;
for( int i = 0 ; i < n ; i ++ ){</pre>
                                                                    origin_u = u;
  for( int j = 0 ; j < n ; j ++ )
  dst[ i ][ j ] = INF;
dst[ i ][ i ] = 0;</pre>
                                                                   void solve() {
                                                                    dfs( 1, 1 );
while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
                                                                    sort( que, que + q, [](const Que& x, const Que& y) {
  }
                                                                     return tie( block_id[ x.u ], dfn[ x.v ] )
 void add_edge( int ui , int vi , int wi ){
                                                                          < tie( block_id[ y.u ], dfn[ y.v ] );
 dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
                                                                   } );
```

```
int U = 1, V = 1;
for ( int i = 0 ; i < q ; ++ i ) {
 pass( U, que[ i ].u );
 pass( V, que[ i ].v );
  // we could get our answer of que[ i ].id
Method 2:
dfs u:
push u
iterate subtree
push u
Let P = LCA(u, v), and St(u) <= St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
3.15 Tree Hashing
uint64_t hsah(int u, int f) {
uint64_t r = 127;
for (int v : G[ u ]) if (v != f) {
 uint64_t hh = hsah(v, u);
 r=(r+(hh*hh)%1010101333)%1011820613;
return r;
3.16 Virtural Tree
inline bool cmp(const int &i, const int &j) {
return dfn[i] < dfn[j];</pre>
void build(int vectrices[], int k) {
static int stk[MAX_N];
sort(vectrices, vectrices + k, cmp);
stk[sz++] = 0;
for (int i = 0; i < k; ++i) {
 int u = vectrices[i], lca = LCA(u, stk[sz - 1]);
  if (lca == stk[sz - 1]) stk[sz++] = u;
 else {
  while (sz \ge 2 \&\& dep[stk[sz - 2]] \ge dep[lca]) {
   addEdge(stk[sz - 2], stk[sz - 1]);
   sz--:
  if (stk[sz - 1] != lca) {
   addEdge(lca, stk[--sz]);
   stk[sz++] = lca, vectrices[cnt++] = lca;
   stk[sz++] = u;
for (int i = 0; i < sz - 1; ++i)
 addEdge(stk[i], stk[i + 1]);
    Matching & Flow
4
4.1 Bipartite Matching
struct BipartiteMatching {
vector<int> X[N];
int fX[N], fY[N], n;
```

```
bitset<N> vis;
bool dfs(int x)
 for (auto i:X[x]) {
  if (vis[i]) continue;
  vis[i] = true;
  if (fY[i]==-1 || dfs(fY[i])){
   fY[fX[x] = i] = x;
   return true;
  }
 }
 return false;
void init(int n_, int m) {
 vis.reset();
fill(X, X + (n = n_), vector<int>());
memset(fX, -1, sizeof(int) * n);
 memset(fY, -1, sizeof(int) * m);
void add_edge(int x, int y){
X[x].push_back(y); }
```

```
int solve() { // return how many pair matched
  int cnt = 0;
  for(int i=0;i<n;i++) {</pre>
   vis.reset();
   cnt += dfs(i):
  return cnt;
};
4.2 Dijkstra Cost Flow
// kN = #(vertices)
// MCMF.{Init, AddEdge, MincostMaxflow}
// MincostMaxflow(source, sink, flow_limit, &cost)
// => flow
using Pii = pair<int, int>
constexpr int kInf = 0x3f3f3f3f, kN = 500;
struct Edge {
int to, rev, cost, flow;
};
struct MCMF { // 0-based
int n{}, m{}, s{}, t{};
 vector<Edge> graph[kN];
 // Larger range for relabeling
 int64_t dis[kN] = {}, h[kN] = {};
 int p[kN] = {};
 void Init(int nn) {
 n = nn;
  for (int i = 0; i < n; i++) graph[i].clear();</pre>
 void AddEdge(int u, int v, int f, int c) {
  graph[u].push_back({v
   static_cast<int>(graph[v].size()), c, f});
  graph[v].push_back(
   {u, static_cast<int>(graph[u].size()) - 1,
    -c, 0});
 bool Dijkstra(int &max_flow, int64_t &cost) {
  priority_queue<Pii, vector<Pii>, greater<>> pq;
  fill_n(dis, n, kInf);
  dis[s] = 0;
  pq.emplace(0, s);
  while (!pq.empty()) {
   auto u = pq.top();
   pq.pop();
   int v = u.second;
   if (dis[v] < u.first) continue;</pre>
   for (auto &e : graph[v]) {
    auto new_dis =
     dis[v] + e.cost + h[v] - h[e.to];
    if (e.flow > 0 && dis[e.to] > new_dis) {
     dis[e.to] = new_dis;
     p[e.to] = e.rev
     pq.emplace(dis[e.to], e.to);
  if (dis[t] == kInf) return false;
  for (int i = 0; i < n; i++) h[i] += dis[i];
  int d = max_flow;
  for (int u = t; u != s;
u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   d = min(d, graph[e.to][e.rev].flow);
  max_flow -= d;
  cost += int64_t(d) * h[t];
  for (int u = t; u != s;
    u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   e.flow += d;
   graph[e.to][e.rev].flow -= d;
  return true:
 int MincostMaxflow(
  int ss, int tt, int max_flow, int64_t &cost) {
  this->s = ss, this->t = tt;
  cost = 0;
  fill_n(h, n, 0);
  auto orig_max_flow = max_flow;
  while (Dijkstra(max_flow, cost) && max_flow) {}
```

return orig_max_flow - max_flow;

```
};
4.3 Dinic
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
     int to, rev;
     Cap cap;
  }:
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
     lv.assign(n, -1);
     queue<int> bfs;
     bfs.push(st); lv[st] = 0;
     while (not bfs.empty()){
       int u = bfs.front(); bfs.pop();
       for (auto e: G[u]) {
         if (e.cap <= 0 or lv[e.to]!=-1) continue;
bfs.push(e.to); lv[e.to] = lv[u] + 1;
     }
     return lv[ed] != -1;
  Cap DFS(int u, Cap f){
     if (u == ed) return f;
     Cap ret = 0:
     for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
       auto &e = G[u][i];
       if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
       Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
       if (f == 0) return ret;
     if (ret == 0) lv[u] = -1;
     return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
  G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
     st = st_, ed = ed_; Cap ret = 0;
     while (BFS()) {
       idx.assign(n, 0);
       Cap f = DFS(st, numeric_limits<Cap>::max());
       ret += f;
       if (f == 0) break;
     return ret;
  }
};
      Flow Models
```

- · Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect v
 ightarrow T with capacity -in(v).
 - To maximize, connect t o s with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.

 To minimize, let f be the maximum flow from S to T. Connect
 - t o s with capacity ∞ and let the flow from S to T be f'. If $f+f'\neq \sum_{v\in V, in(v)>0}in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e\,+\,f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching ${\cal M}$ on bipartite graph(X, Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X
 - 3. $x \in X$ is chosen iff x is unvisited. 4. $y \in Y$ is chosen iff y is visited.

- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect $x \to y$ with (cost,cap) = (c,1) if c>0, otherwise connect $y\to x$ with (cost, cap)=(-c,1)
 - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v) > 0, connect $S \to v$ with (cost, cap) =(0, d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =(0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \rightarrow v$, $v \in G$ with capacity K
 - 4. For each edge (u, v, w) in G, connect $u \to v$ and $v \to u$ with capacity
 - 5. For $v \in {\it G}$, connect it with sink $v \to t$ with capacity K + 2T $(\sum_{e \in E(v)} w(e)) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v
 - 3. Find the minimum weight perfect matching on G^{\prime} .
- · Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge
 - (v,t) with capacity $-p_v$. 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

namespace matching {

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with ca-
- pacity c_y . 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

General Graph Matching

```
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
g[u].push_back(v);
 g[v].push_back(u);
int Find(int u) {
 return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
 static int tk = 0; tk++;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
  }
 }
void Blossom(int x, int y, int 1) {
  while (Find(x) != 1) {
  pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
  if (fa[x] == x) fa[x] = 1;
  if (fa[y] == y) fa[y] = 1;
  x = pre[y];
bool Bfs(int r, int n) {
 for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;
 while (!q.empty()) q.pop();
 q.push(r);
 s[r] = 0;
```

for(int j=i+1;j<=n;++j){</pre>

```
while (!q.empty()) {
                                                                  if(g[j]==t && flow.connect(j))g[j]=i; // check if i
  int x = q.front(); q.pop();
                                                                   can reach i
  for (int u : g[x]) {
  if (s[u] == -1) {
                                                                 }
    pre[u] = x, s[u] = 1;
                                                                return rt:
    if (match[u] == n) {
                                                              }
     for (int a = u, b = x, last; b != n; a = last, b =
                                                               4.8 Kuhn Munkres
     pre[a])
      last = match[b], match[b] = a, match[a] = b;
     return true;
                                                              private:
                                                                static constexpr 11d INF = 1LL << 60;</pre>
                                                                vector<lld> hl,hr,slk;
    q.push(match[u]);
                                                                vector<int> fl,fr,pre,qu;
    s[match[u]] = 0;
   } else if (!s[u] && Find(u) != Find(x)) {
                                                                vector<vector<lld>> w;
    int 1 = LCA(u, x, n);
                                                                vector<bool> v1,vr;
    Blossom(x, u, 1);
                                                                int n, ql, qr;
                                                               bool check(int x) {
  if (v1[x] = true, f1[x] != -1)
    Blossom(u, x, 1);
  }
                                                                  return vr[qu[qr++] = f1[x]] = true;
                                                                 while (x != -1) swap(x, fr[fl[x] = pre[x]]);
 return false;
                                                                 return false;
int Solve(int n) {
                                                                void bfs(int s) {
                                                                 fill(slk.begin(), slk.end(), INF);
 int res = 0:
 for (int x = 0; x < n; ++x) {
                                                                 fill(vl.begin(), vl.end(), false);
  if (match[x] == n) res += Bfs(x, n);
                                                                 fill(vr.begin(), vr.end(), false);
                                                                 q1 = qr = 0;
                                                                 vr[qu[qr++] = s] = true;
 return res;
}}
                                                                 while (true) {
                                                                  11d d;
                                                                  while (ql < qr) {</pre>
4.6 Global Min-Cut
                                                                   for (int x = 0, y = qu[ql++]; x < n; ++x) {
const int maxn = 500 + 5;
                                                                    if(!v1[x]&&slk[x]>=(d=h1[x]+hr[y]-w[x][y])){
int w[maxn][maxn], g[maxn];
                                                                    if (pre[x] = y, d) slk[x] = d;
bool v[maxn], del[maxn];
                                                                     else if (!check(x)) return;
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
                                                                   }
pair<int, int> phase(int n) {
                                                                  d = INF;
 memset(v, false, sizeof(v));
                                                                  for (int x = 0; x < n; ++x)
 memset(g, 0, sizeof(g));
                                                                   if (!v1[x] \&\& d > s1k[x]) d = s1k[x];
 int s = -1, t = -1;
                                                                  for (int x = 0; x < n; ++x) {
  if (v1[x]) h1[x] += d;
 while (true) {
  int c = -1;
                                                                   else slk[x] -= d;
  for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
                                                                   if (vr[x]) hr[x] -= d;
   if (c == -1 \mid \mid g[i] > g[c]) c = i;
                                                                  for (int x = 0; x < n; ++x)
                                                                   if (!vl[x] && !slk[x] && !check(x)) return;
  if (c == -1) break;
  v[s = t, t = c] = true;
  for (int i = 0; i < n; ++i) {
                                                              public:
   if (del[i] || v[i]) continue;
                                                                void init( int n_ ) {
   g[i] += w[c][i];
                                                                 qu.resize(n = n_);
  }
                                                                fl.assign(n, -1); fr.assign(n, -1);
                                                                hr.assign(n, 0); hl.resize(n);
w.assign(n, vector<lld>(n));
 return make_pair(s, t);
                                                                slk.resize(n); pre.resize(n);
int mincut(int n) {
                                                                vl.resize(n); vr.resize(n);
 int cut = 1e9;
 memset(del, false, sizeof(del));
                                                                void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 for (int i = 0; i < n - 1; ++i) {
                                                                11d solve() {
  int s, t; tie(s, t) = phase(n);
                                                                 for (int i = 0; i < n; ++i)</pre>
  del[t] = true; cut = min(cut, g[t]);
                                                                 hl[i] = *max_element(w[i].begin(), w[i].end());
  for (int j = 0; j < n; ++j) {
                                                                 for (int i = 0; i < n; ++i) bfs(i);</pre>
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
                                                                11d res = 0;
                                                                for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
                                                                 return res;
return cut;
                                                               }
                                                              } km;
4.7 GomoryHu Tree
                                                               4.9 Minimum Cost Circulation
                                                              struct Edge { int to, cap, rev, cost; };
int g[maxn];
vector<edge> GomoryHu(int n){
                                                              vector<Edge> g[kN];
 vector<edge> rt:
                                                              int dist[kN], pv[kN], ed[kN];
 for(int i=1;i<=n;++i)g[i]=1;</pre>
                                                              bool mark[kN];
 for(int i=2;i<=n;++i){</pre>
                                                              int NegativeCycle(int n) {
                                                               memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
  int t=q[i]
  flow.reset(); // clear flows on all edge
  rt.push_back({i,t,flow(i,t)});
                                                                int upd = -1;
                                                                for (int i = 0; i <= n; ++i)
  flow.walk(i); // bfs points that connected to i (use
    edges not fully flow)
                                                                for (int j = 0; j < n; ++j) {
```

int idx = 0;

```
for (auto &e : g[j]) {
                                                              Cap mw=INF_CAP;
    if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
                                                              for(int i=edd;i!=ori;i=fa[i])
     dist[e.to] = dist[j] + e.cost;
                                                              mw=min(mw,G[fa[i]][wh[i]].cap);
                                                              for (int i=edd;i!=ori;i=fa[i]){
     pv[e.to] = j, ed[e.to] = idx;
     if (i == n) {
                                                               auto &eg=G[fa[i]][wh[i]];
      upd = j;
                                                               eg.cap -= mw;
      while(!mark[upd])mark[upd]=1,upd=pv[upd];
                                                               G[eg.to][eg.back].cap+=mw;
      return upd;
                                                              return {mw, dis[edd]};
                                                           public:
    idx++;
                                                             void init(int n){
 }
                                                             G.clear();G.resize(n);
                                                              fa.resize(n);wh.resize(n);
return -1;
                                                              inq.resize(n); dis.resize(n);
int Solve(int n) {
                                                             void add_edge(int st, int ed, Cap c, Wei w){
                                                             G[st].emplace_back(ed,SZ(G[ed]),c,w);
int rt = -1, ans = 0;
while ((rt = NegativeCycle(n)) >= 0) {
                                                             G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
 memset(mark, false, sizeof(mark));
 vector<pair<int, int>> cyc;
                                                             PCW solve(int a, int b){
 while (!mark[rt]) {
                                                              ori = a, edd = b;
                                                              Cap cc=0; Wei ww=0;
  cyc.emplace_back(pv[rt], ed[rt]);
  mark[rt] = true;
                                                              while(true)
  rt = pv[rt];
                                                              PCW ret=SPFA();
                                                               if(ret.first==-1) break;
 reverse(cyc.begin(), cyc.end());
                                                              cc+=ret.first;
 int cap = kInf;
                                                               ww+=ret.first * ret.second;
 for (auto &i : cyc)
                                                              }
  auto &e = g[i.first][i.second];
                                                              return {cc,ww};
  cap = min(cap, e.cap);
                                                           } mcmf;
 for (auto &i : cyc)
                                                            4.11
                                                                  Minimum Weight Matching (Clique version)
  auto &e = g[i.first][i.second];
  e.cap -= cap;
                                                           struct Graph {
  g[e.to][e.rev].cap += cap;
                                                             // 0-base (Perfect Match)
  ans += e.cost * cap;
                                                             int n, edge[MXN][MXN];
 }
                                                             int match[MXN], dis[MXN], onstk[MXN];
                                                             vector<int> stk;
return ans;
                                                             void init(int _n) {
                                                             n = _n;
                                                             for (int i=0; i<n; i++) for (int j=0; j<n; j++)</pre>
4.10 Minimum Cost Maximum Flow
                                                               edge[i][j] = 0;
class MiniCostMaxiFlow{
using Cap = int; using Wei = int64_t;
                                                             void set_edge(int u, int v, int w) {
using PCW = pair<Cap,Wei>;
                                                              edge[u][v] = edge[v][u] = w; }
                                                             bool SPFA(int u){
static constexpr Cap INF_CAP = 1 << 30;</pre>
static constexpr Wei INF_WEI = 1LL<<60;</pre>
                                                              if (onstk[u]) return true;
                                                              stk.PB(u); onstk[u] = 1;
private:
struct Edge{
                                                              for (int v=0; v<n; v++){</pre>
 int to, back;
                                                               if (u != v && match[u] != v && !onstk[v]){
 Cap cap; Wei wei;
                                                                int m = match[v]
                                                                if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
 Edge() {}
 Edge(int a,int b, Cap c, Wei d):
                                                                 dis[m] = dis[u] - edge[v][m] + edge[u][v];
  to(a),back(b),cap(c),wei(d) {}
                                                                 onstk[v] = 1;
                                                                 stk.PB(v)
int ori, edd;
                                                                 if (SPFA(m)) return true;
vector<vector<Edge>> G;
                                                                 stk.pop_back();
vector<int> fa, wh;
                                                                 onstk[v] = 0;
vector<bool> inq;
vector<Wei> dis;
PCW SPFA(){
                                                             onstk[u] = 0; stk.pop_back();
 fill(inq.begin(),inq.end(),false);
                                                              return false;
 fill(dis.begin(), dis.end(), INF_WEI);
  queue<int> qq; qq.push(ori);
 dis[ori] = 0;
                                                             int solve() { // find a match
 while(not qq.empty()){
                                                             for (int i=0; i<n; i+=2){
   int u=qq.front();qq.pop();
                                                               match[i] = i+1;
  inq[u] = false
                                                              match[i+1] = i;
   for(int i=0;i<SZ(G[u]);++i){</pre>
   Edge e=G[u][i];
                                                              while (true){
                                                               int found = 0;
    int v=e.to; Wei d=e.wei;
    if(e.cap<=0||dis[v]<=dis[u]+d)</pre>
                                                               for (int i=0; i<n; i++)</pre>
     continue;
                                                                dis[i] = onstk[i] = 0;
    dis[v] = dis[u] + d;
                                                               for (int i=0; i<n; i++){</pre>
                                                                stk.clear()
    fa[v] = u, wh[v] = i;
   if (inq[v]) continue;
                                                                if (!onstk[i] && SPFA(i)){
    qq.push(v);
                                                                 found = 1
    inq[v] = true;
                                                                 while (SZ(stk)>=2){
                                                                  int u = stk.back(); stk.pop_back();
                                                                  int v = stk.back(); stk.pop_back();
  if(dis[edd]==INF_WEI) return {-1, -1};
                                                                 match[u] = v;
```

```
match[v] = u;
    }
    }
    if (!found) break;
}
    int ret = 0;
    for (int i=0; i<n; i++)
        ret += edge[i][match[i]];
    return ret>>1;
}
} graph;
```

5 Math

5.1 $\lfloor \frac{n}{i} \rfloor$ Enumeration

$$T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor$$

5.2 Strling Number

5.2.1 First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1) \dots (x+n-1) = \sum_{k=0}^n S_1(n,k) x^k$$

$$g(x) = x(x+1) \dots (x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.2.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b==(x/y)*a;
}
```

5.4 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
vector<T> d(output.size() + 1), me, he;
for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
 for (size_t j = 0; j < me.size(); ++j)
  d[i] += output[i - j - 2] * me[j];
if ((d[i] -= output[i - 1]) == 0) continue;</pre>
  if (me.empty()) {
   me.resize(f = i);
   continue:
  }
 vector<T> o(i - f - 1);
T k = -d[i] / d[f]; o.push_back(-k);
  for (T x : he) o.push_back(x * k);
  if (o.size() < me.size()) o.resize(me.size());</pre>
  for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
  if (i-f+he.size() >= me.size()) he = me, f = i;
  me = o;
return me;
```

5.5 Charateristic Polynomial

```
vector<vector<int>> Hessenberg(const vector<vector<int
    >> &A) {
 int N = A.size();
 vector<vector<int>> H = A;
 for (int i = 0; i < N - 2; ++i) {
  if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {
    if (H[j][i]) {
      for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
    ][k]);
      for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k]
    ][j]);
     break;
    }
   }
  if (!H[i + 1][i]) continue;
int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
   int coef = 1LL * val * H[j][i] % kP;
for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL</pre>
      * H[i + 1][k] * (kP - coef)) % kP;
   for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
 return H;
vector<int> CharacteristicPoly(const vector<vector<int</pre>
    >> &A) {
 int N = A.size();
 auto H = Hessenberg(A);
 for (int i = 0; i < N; ++i) {
  for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
 vector<vector<int>>> P(N + 1, vector<int>(N + 1));
 P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
  P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
    11:
  int val = 1;
  for (int j = i - 1; j >= 0; --j) {
  int coef = 1LL * val * H[j][i - 1] % kP;
   for (int k = 0; k \le j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 if (N & 1) {
  for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
 return P[N];
```

5.6 Chinese Remainder

```
1ld crt(lld ans[], lld pri[], int n){
 lld M = 1, ret = 0;
 for(int i=0;i<n;i++) M *= pri[i];</pre>
 for(int i=0;i<n;i++){</pre>
  1ld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
  ret += (ans[i]*(M/pri[i])%M * iv)%M;
  ret %= M;
 return ret:
}
/*
Another:
x = a1 \% m1
x = a2 \% m2
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
0 <= x < lcm(m1, m2)
```

5.7 De-Bruijn

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
```

```
if (t > n) {
 if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
} else {
 aux[t] = aux[t - p];
 db(t + 1, p, n, k);
for (int i = aux[t - p] + 1; i < k; ++i) {
   aux[t] = i;
   db(t + 1, t, n, k);
int de_bruijn(int k, int n) {
// return cyclic string of len k^n s.t. every string
// of len n using k char appears as a substring.
if (k == 1) {
 res[0] = 0;
 return 1;
for (int i = 0; i < k * n; i++) aux[i] = 0;
sz = 0;
db(1, 1, n, k);
return sz;
```

5.8 DiscreteLog

```
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
  // x^? \equiv y (mod M)
Int t = 1, c = 0, g = 1;
  for (Int M<sub>_</sub> = M; M<sub>_</sub> > 0; M<sub>_</sub> >>= 1)
    g = g * x % M;
  for (g = gcd(g, M); t % g != 0; ++c) {
    if (t == y) return c;
t = t * x % M;
  if (y % g != 0) return -1;
  t /= g, y /= g, M /= g;
 Int h = 0, gs = 1;
for (; h * h < M; ++h) gs = gs * x % M;
  unordered_map<Int, Int> bs;
  for (Int s = 0; s < h; bs[y] = ++s)
    y = y * x % M;
  for (Int s = 0; s < M; s += h) {
    t = t * gs % M;
    if (bs.count(t)) return c + s + h - bs[t];
  return -1:
```

5.9 Extended Euler

$$a^b \equiv \begin{cases} a^b \mod \varphi(m) + \varphi(m) & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m$$

5.10 ExtendedFloorSum

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \\ -\frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \\ h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ -2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

5.11 Fast Fourier Transform

```
const int mod = 1000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
  static_assert (M1 <= M2 && M2 <= M3);</pre>
  constexpr int64_t r12 = modpow(M1, M2-2, M2);
  constexpr int64_t r13 = modpow(M1, M3-2, M3);
constexpr int64_t r23 = modpow(M2, M3-2, M3);
  constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
  B = (B - A + M2) * r12 % M2;

C = (C - A + M3) * r13 % M3;
  C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
 for (int i = 0; i <= maxn; i++)</pre>
  omega[i] = cplx(cos(2 * pi * j / maxn),
     sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
 int z = __builtin_ctz(n) - 1;
 for (int i = 0; i < n; ++i) {</pre>
  int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^{=(i >> j & 1) << (z - j);
  if (x > i) swap(v[x], v[i]);
 for (int s = 2; s <= n; s <<= 1) {
  int z = s \gg 1;
  for (int i = 0; i < n; i += s) {
   for (int k = 0; k < z; ++k) {
  cplx x = v[i + z + k] * omega[maxn / s * k];
    v[i + z + k] = v[i + k] - x;
    v[i+k] = v[i+k] + x;
void ifft(vector<cplx> &v, int n) {
 fft(v, n); reverse(v.begin() + 1, v.end());
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
VL convolution(const VI &a, const VI &b) {
 // Should be able to handle N <= 10^5, C <= 10^4
 int sz = 1:
 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 vector<cplx> v(sz);
 for (int i = 0; i < sz; ++i) {
  double re = i < a.size() ? a[i] : 0;</pre>
  double im = i < b.size() ? b[i] : 0;</pre>
  v[i] = cplx(re, im);
 fft(v, sz);
 for (int i = 0; i <= sz / 2; ++i) {
  int j = (sz - i) & (sz - 1);
  cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
     * cplx(0, -0.25);
  if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
     ].conj()) * cplx(0, -0.25);
  v[i] = x;
 ifft(v, sz);
 VL c(sz);
 for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
VI convolution_mod(const VI &a, const VI &b, int p) {
 while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
 vector<cplx> fa(sz), fb(sz);
 for (int i = 0; i < (int)a.size(); ++i)</pre>
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 for (int i = 0; i < (int)b.size(); ++i)</pre>
 fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
```

```
fft(fa, sz), fft(fb, sz);
                                                                  int d2 = d << 1;
                                                                  for( int s = 0 ; s < N ; s += d2 )
 double r = 0.25 / sz;
                                                                  for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
    LL ta = x[ i ] , tb = x[ j ];
    x[ i ] = ta+tb;
cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);

for (int i = 0; i <= (sz >> 1); ++i) {

  int j = (sz - i) & (sz - 1);
  cplx a1 = (fa[i] + fa[j].conj());
                                                                    x[ j ] = ta-tb;
                                                                   if(x[i] >= MOD ) x[i] -= MOD;
if(x[j] < 0 ) x[j] += MOD;</pre>
  cplx a2 = (fa[i] - fa[j].conj()) * r2;
  cplx b1 = (fb[i] + fb[j].conj()) * r3;
  cplx b2 = (fb[i] - fb[j].conj()) * r4;
  if (i != j) {
   cplx c1 = (fa[j] + fa[i].conj());
                                                                if( inv )
   cplx c2 = (fa[j] - fa[i].conj()) * r2;
                                                                 for( int i = 0 ; i < N ; i++ ) {</pre>
                                                                  x[i] *= inv(N, MOD);
x[i] %= MOD;
   cplx d1 = (fb[j] + fb[i].conj()) * r3;
   cplx d2 = (fb[j] - fb[i].conj()) * r4;
   fa[i] = c1 * d1 + c2 * d2 * r5;
   fb[i] = c1 * d2 + c2 * d1;
                                                               5.14
                                                                       Gauss Elimination
  fa[j] = a1 * b1 + a2 * b2 * r5;
  fb[j] = a1 * b2 + a2 * b1;
                                                               void gauss(vector<vector<double>> &d) {
                                                                int n = d.size(), m = d[0].size();
 fft(fa, sz), fft(fb, sz);
                                                                for (int i = 0; i < m; ++i) {
 vector<int> res(sz);
                                                                  int p = -1;
 for (int i = 0; i < sz; ++i) {
                                                                  for (int j = i; j < n; ++j) {</pre>
 long long a = round(fa[i].re), b = round(fb[i].re),
                                                                   if (fabs(d[j][i]) < eps) continue;</pre>
       c = round(fa[i].im);
                                                                   if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
  res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
                                                                  if (p == -1) continue;
 return res;
                                                                 for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]); for (int j = 0; j < n; ++j) {
}}
                                                                   if (i == j) continue;
5.12 FloorSum
                                                                   double z = d[j][i] / d[i][i];
// @param n \cdot n < 2^32
                                                                   for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
// @param m `1 <= m < 2^32`
// @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
 llu ans = 0;
 while (true) -
                                                               5.15 Miller Rabin
  if (a >= m) {
                                                               bool isprime(llu x){
   ans += n * (n - 1) / 2 * (a / m); a %= m;
                                                                static llu magic[]={2,325,9375,28178.\
                                                                          450775,9780504,1795265022};
  if (b >= m) {
                                                                static auto witn=[](llu a,llu u,llu n,int t)
   ans += n * (b / m); b %= m;
                                                                ->bool{
                                                                  if (!(a = mpow(a%n,u,n)))return 0;
  llu y_max = a * n + b;
                                                                  while(t--){
  if (y_max < m) break;</pre>
                                                                  llu a2=mul(a,a,n);
  // y_max < m * (n + 1)
                                                                   if(a2==1 && a!=1 && a!=n-1)
 // floor(y_max / m) <= n
                                                                   return 1:
 n = (1lu)(y_max / m), b = (1lu)(y_max % m);
                                                                  a = a2;
  swap(m, a);
                                                                 }
 }
                                                                  return a!=1;
 return ans;
                                                                if(x<2)return 0;</pre>
11d floor_sum(11d n, 11d m, 11d a, 11d b) {
                                                                if(!(x&1))return x==2;
 11u ans = 0;
                                                                llu x1=x-1;int t=0;
 if (a < 0) {
                                                                while(!(x1&1))x1>>=1,t++;
 11u \ a2 = (a \% m + m) \% m;
                                                                for(llu m:magic)if(witn(m,x1,x,t))return 0;
  ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
                                                                return 1:
 a = a2:
 if (b < 0) {
                                                               5.16 NTT
 11u b2 = (b \% m + m) \% m;
  ans -= 1ULL * n * ((b2 - b) / m);
                                                               template <int mod, int G, int maxn>
                                                               struct NTT {
                                                                static_assert (maxn == (maxn & -maxn));
 return ans + floor_sum_unsigned(n, m, a, b);
                                                                int roots[maxn];
                                                                NTT () {
                                                                  int r = modpow(G, (mod - 1) / maxn);
5.13 FWT
                                                                  for (int i = maxn >> 1; i; i >>= 1) {
                                                                  roots[i] = 1;
/* xor convolution:
* x = (x0,x1) , y = (y0,y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                   for (int j = 1; j < i; j++)
                                                                   roots[i + j] = modmul(roots[i + j - 1], r);
                                                                   r = modmul(r, r);
 * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
 * z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
* z = (1/2) * z''
                                                                // n must be 2^k, and 0 \le F[i] < mod
 * or convolution:
                                                                void operator()(int F[], int n, bool inv = false) {
                                                                 for (int i = 0, j = 0; i < n; i++) {
  if (i < j) swap(F[i], F[j]);
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
 * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
                                                                   for (int k = n > 1; (j^k < k; k > = 1);
const LL MOD = 1e9+7;
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
                                                                  for (int s = 1; s < n; s *= 2) {
for( int d = 1 ; d < N ; d <<= 1 ) {
                                                                  for (int i = 0; i < n; i += s * 2) {
```

```
for (int j = 0; j < s; j++) {
                                                                return ret;
     int a = F[i+j];
     int b = modmul(F[i+j+s], roots[s+j]);
                                                               11d pi_count(11d);
     F[i+j] = modadd(a, b); // a + b

F[i+j+s] = modsub(a, b); // a - b
                                                               11d P2(11d m, 11d n) {
                                                                11d sm = square_root(m), ret = 0;
                                                                for(lld i = n+1;primes[i]<=sm;i++)</pre>
   }
                                                                 ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
                                                                return ret:
  if (inv) {
   int invn = modinv(n);
                                                               11d pi_count(11d m) {
   for (int i = 0; i < n; i++)</pre>
                                                                if(m < N) return pi[m];</pre>
    F[i] = modmul(F[i], invn);
                                                                11d n = pi_count(cube_root(m));
   reverse(F + 1, F + n);
                                                                return phi(m, n) + n - 1 - P2(m, n);
                                                               5.20 Pollard Rho
NTT<2013265921, 31, 1048576> ntt;
                                                               // does not work when n is prime
                                                               // return any non-trivial factor
5.17
      Range Sieve
                                                               llu pollard_rho(llu n){
                                                                static auto f=[](llu x,llu k,llu m){
const int MAX_SQRT_B = 50000;
                                                                 return add(k,mul(x,x,m),m);};
const int MAX_L = 200000 + 5;
                                                                if (!(n&1)) return 2;
bool is_prime_small[MAX_SQRT_B], is_prime[MAX_L];
                                                                mt19937 rnd(120821011);
void sieve(lld 1, lld r){ // [1, r)
                                                                while(true){
 for(lld i=2;i*i<r;i++) is_prime_small[i] = true;</pre>
                                                                 llu y=2,yy=y,x=rnd()%n,t=1;
 for(lld i=1;i<r;i++) is_prime[i-1] = true;
if(l==1) is_prime[0] = false;</pre>
                                                                 for(llu sz=2;t==1;sz<<=1,y=yy) {</pre>
                                                                  for(llu i=0;t==1&&i<sz;++i){</pre>
 for(lld i=2;i*i<r;i++){</pre>
                                                                   yy=f(yy,x,n);
  if(!is_prime_small[i]) continue;
                                                                   t=gcd(yy>y?yy-y:y-yy,n);
  for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;</pre>
  for(1ld j=std::max(2LL, (1+i-1)/i)*i;j<r;j+=i)</pre>
   is_prime[j-1]=false;
                                                                 if(t!=1&&t!=n) return t;
}
                                                               }
5.18
     Partition Number
                                                                     Polynomial Operations
int b = sqrt(n);
                                                               using V = vector<int>
ans[0] = tmp[0] = 1;
                                                               #define fi(s, n) for (int i = int(s); i < int(n); ++i)
for (int i = 1; i <= b; i++) {
                                                               template <int mod, int G, int maxn> struct Poly : V {
 for (int rep = 0; rep < 2; rep++)
                                                                static uint32_t n2k(uint32_t n) {
  for (int j = i; j <= n - i * i; j++)
                                                                 if (n <= 1) return 1;
   modadd(tmp[j], tmp[j-i]);
                                                                 return 1u << (32 - __builtin_clz(n - 1));</pre>
 for (int j = i * i; j <= n; j++)
  modadd(ans[j], tmp[j - i * i]);
                                                                static NTT<mod,G,maxn> ntt; // coefficients in [0, P)
                                                                explicit Poly(int n = 1) : V(n) {}
                                                                Poly(const V &v) : V(v) {}
5.19 Pi Count (Linear Sieve)
                                                                Poly(const Poly &p, size_t n) : V(n) {
static constexpr int N = 1000000 + 5;
                                                                 copy_n(p.data(), min(p.size(), n), data());
lld pi[N];
                                                                Poly &irev() { return reverse(data(), data() + size())
vector<int> primes;
bool sieved[N];
                                                                     *this; }
1ld cube_root(lld x){
                                                                Poly &isz(int _n) { return resize(_n), *this;
                                                                Poly &iadd(const Poly &rhs) { // n() == rhs.n()
1ld s=cbrt(x-static_cast<long double>(0.1));
                                                                 fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
 while(s*s*s <= x) ++s;
 return s-1;
                                                                 return *this;
11d square_root(11d x){
                                                                Poly &imul(int k) {
lld s=sqrt(x-static_cast<long double>(0.1));
                                                                 fi(0, size())(*this)[i] = modmul((*this)[i], k);
 while(s*s <= x) ++s;</pre>
                                                                 return *this;
 return s-1;
                                                                Poly Mul(const Poly &rhs) const {
void init(){
                                                                 const int _n = n2k(size() + rhs.size() - 1);
                                                                 Poly X(*this, _n), Y(rhs, _n);
primes.reserve(N);
                                                                 ntt(X.data(), _n), ntt(Y.data(), _n)
fi(0, _n) X[i] = modmul(X[i], Y[i]);
 primes.push_back(1);
 for(int i=2;i<N;i++) {</pre>
 if(!sieved[i]) primes.push_back(i);
                                                                 ntt(X.data(), _n, true);
  pi[i] = !sieved[i] + pi[i-1];
                                                                 return X.isz(size() + rhs.size() - 1);
  for(int p: primes) if(p > 1) {
  if(p * i >= N) break;
                                                                Poly Inv() const { // coef[0] != 0
   sieved[p * i] = true;
                                                                 if (size() == 1) return V{modinv(*begin())};
                                                                 const int _n = n2k(size() * 2);
Poly Xi = Poly(*this, (size() + 1) / 2).Inv().isz(_n)
   if(p % i == 0) break;
                                                                 , Y(*this, _n);
ntt(Xi.data(), _n), ntt(Y.data(), _n);
fi(0, _n) Xi[i] = modmul(Xi[i], modsub(2, modmul(Xi[i]))
11d phi(lld m, lld n) {
 static constexpr int MM = 80000, NN = 500;
                                                                    ], Y[i])));
 static lld val[MM][NN];
                                                                 ntt(Xi.data(), _n, true);
                                                                 return Xi.isz(size());
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
 if(n == 0) return m;
                                                                Poly Sqrt() const { // coef[0] \in [1, mod)^2
 if(primes[n] >= m) return 1;
 1ld ret = phi(m,n-1)-phi(m/primes[n],n-1);
                                                                 if (size() == 1) return V{QuadraticResidue((*this)
 if(m < MM\&n < NN) val[m][n] = ret+1;
                                                                    [0], mod)};
```

```
Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
                                                              nk2 = nk2 * 10 + c - '0';
                                                              if (nk2 * nz >= size())
   size()):
                                                               return Poly(size());
 return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
                                                              nk2 \% = mod - 1;
                                                             if (!nk && !nk2) return Poly(V{1}, size());
pair<Poly, Poly> DivMod(const Poly &rhs) const {
if (size() < rhs.size()) return {V{0}, *this};</pre>
                                                             Poly X = V(data() + nz, data() + size() - nz * (nk2 -
 const int _n = size() - rhs.size() + 1;
                                                                1));
Poly X(rhs); X.irev().isz(_n);
                                                             int x0 = X[0];
Poly Y(*this); Y.irev().isz(_n);
                                                             return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
Poly Q = Y.Mul(X.Inv()).isz(_n).irev();
                                                               modpow(x0, nk2)).irev().isz(size()).irev();
 X = rhs.Mul(Q), Y = *this;
 fi(0, size()) Y[i] = modsub(Y[i], X[i]);
                                                            Poly InvMod(int L) { // (to evaluate linear recursion)
                                                                            // *this * R mod x^L = 1 (*this[0] ==
 return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
                                                             Poly R{1, 0};
Poly Dx() const {
                                                             for (int level = 0; (1 << level) < L; ++level) {</pre>
Poly ret(size() - 1);
                                                              Poly 0 = R.Mul(Poly(data(), min<int>(2 << level,
                                                               size())));
fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
                                                              Poly Q(2 << level); Q[0] = 1;
   1]);
 return ret.isz(max<int>(1, ret.size()));
                                                              for (int j = (1 << level); j < (2 << level); ++j)</pre>
                                                               Q[j] = modsub(mod, O[j]);
Poly Sx() const {
                                                              R = R.Mul(Q).isz(4 << level);
Poly ret(size() + 1);
fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
                                                             return R.isz(L);
   this)[i]);
                                                            static int LinearRecursion(const V &a, const V &c,
 return ret:
                                                               int64_t n) {
Poly _tmul(int nn, const Poly &rhs) const {
                                                             // a_n = \sum c_j a_(n-j)
Poly Y = Mul(rhs).isz(size() + nn - 1);
                                                             const int k = (int)a.size();
 return V(Y.data() + size() - 1, Y.data() + Y.size());
                                                             assert((int)c.size() == k + 1);
                                                             Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
V _eval(const V &x, const vector<Poly> &up) const {
                                                             fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
const int _n = (int)x.size();
                                                             C[k] = 1;
if (!_n) return {};
                                                             while (n) {
                                                              if (n % 2) W = W.Mul(M).DivMod(C).second;
vector<Poly> down(_n * 2);
 // down[1] = DivMod(up[1]).second;
                                                              n /= 2, M = M.Mul(M).DivMod(C).second;
 // fi(2, _n * 2) down[i] = down[i / 2].DivMod(up[i]).
   second;
                                                             int ret = 0;
down[1] = Poly(up[1]).irev().isz(size()).Inv().irev()
                                                             fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
   ._tmul(_n, *this)
                                                             return ret:
 fi(2, _n * 2) down[i] = up[i ^ 1]._tmul(up[i].size()
   - 1, down[i / 2]);
                                                           }:
 V y(_n); fi(0, _n) y[i] = down[_n + i][0];
                                                           #undef fi
 return y;
                                                           using Poly_t = Poly<998244353, 3, 1 << 20>;
                                                           template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
static vector<Poly> _tree1(const V &x) {
                                                           5.22 Quadratic residue
const int _n = (int)x.size();
vector<Poly> up(_n * 2);
                                                           struct S {
 fi(0, _n) up[_n + i] = V\{modsub(mod, x[i]), 1\};
                                                            int MOD, w;
 for(int i=_n-1;i>0;--i) up[i] = up[i * 2].Mul(up[i *
                                                            int64_t x, y;
   2 + 1]);
                                                            S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
 return up;
                                                             : MOD(m), w(w_{-}), x(x_{-}), y(y_{-}) {}
                                                            S operator*(const S &rhs) const {
                                                             int w_ = w;
if (w_ == -1) w_ = rhs.w;
V Eval(const V &x) const { return _eval(x, _tree1(x));
                                                             assert(w_ != -1 and w_ == rhs.w);
static Poly Interpolate(const V &x, const V &y) {
                                                             return { MOD, w_,
 const int _n = (int)x.size();
                                                              (x * rhs.x + y * rhs.y % MOD * w) % MOD,
 vector < Poly > up = _tree1(x), down(_n * 2);
 V z = up[1].Dx()._eval(x, up);
                                                              (x * rhs.y + y * rhs.x) % MOD };
 fi(0, _n) z[i] = modmul(y[i], modinv(z[i]));
 fi(0, n) down[n + i] = V{z[i]};
 for(int i=_n-1;i>0;--i) down[i] = down[i * 2].Mul(up[
                                                           int get_root(int n, int P) {
   i * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
                                                             if (P == 2 \text{ or } n == 0) \text{ return } n;
                                                             if (qpow(n, (P - 1) / 2, P) != 1) return -1;
 return down[1];
                                                             auto check = [&](int x) {
                                                             return qpow(x, (P - 1) / 2, P); };
if (check(n) == P-1) return -1;
Poly Ln() const { // coef[0] == 1
return Dx().Mul(Inv()).Sx().isz(size());
                                                             int64_t a; int w; mt19937 rnd(7122);
                                                             do { a = rnd() % P;
Poly Exp() const \{ // coef[0] == 0 \}
                                                               w = ((a * a' - n) \% P + P) \% P;
 if (size() == 1) return V{1};
Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size)
                                                             } while (check(w) != P - 1);
                                                             return qpow(S(P, w, a, 1), (P + 1) / 2).x;
   ());
Poly Y = X.Ln(); Y[0] = mod - 1;
 fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
                                                           5.23 Simplex
 return X.Mul(Y).isz(size());
                                                           namespace simplex {
                                                           // maximize c^Tx under Ax <= B
Poly Pow(const string &K) const {
                                                           // return VD(n, -inf) if the solution doesn't exist
int nz = 0;
 while (nz < size() && !(*this)[nz]) ++nz;</pre>
                                                           // return VD(n, +inf) if the solution is unbounded
 int nk = 0, nk2 = 0;
                                                           using VD = vector<double>;
 for (char c : K) {
                                                           using VVD = vector<vector<double>>;
 nk = (nk * 10 + c - '0') % mod;
                                                           const double eps = 1e-9;
```

```
const double inf = 1e+9;
                                                                 #define IM imag
                                                                 #define RE real
int n, m;
VVD d;
                                                                 using lld = int64_t;
vector<int> p, q;
void pivot(int r, int s) {
                                                                 using llf = long double;
                                                                 using Point = std::complex<lld>;
                                                                 using Pointf = std::complex<llf>
 double inv = 1.0 / d[r][s];
 for (int i = 0; i < m + 2; ++i)
                                                                 auto toPointf(Point p) { return Pointf(IM(p), RE(p)); }
  for (int j = 0; j < n + 2; ++j)
                                                                 int sgn(lld x) { return (x > 0) - (x < 0); }
   if (i != r && j != s)
                                                                 11d dot(Point a, Point b) { return RE(conj(a) * b); }
    d[i][j] -= d[r][j] * d[i][s] * inv;
                                                                 11d cross(Point a, Point b) { return IM(conj(a) * b); }
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
                                                                 int ori(Point a, Point b, Point c) {
 for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
                                                                  return sgn(cross(b - a, c - a));
 d[r][s] = inv; swap(p[r], q[s]);
                                                                 bool operator<(const Point &a, const Point &b) {</pre>
                                                                  return RE(a) != RE(b) ? RE(a) < RE(b) : IM(a) < IM(b);
bool phase(int z) {
 int x = m + z
 while (true) {
                                                                 int argCmp(Point a, Point b) {
                                                                  // -1 / 0 / 1 <-> < / == / > (atan2)
  int s = -1;
                                                                  int qa = (IM(a) == 0
  for (int i = 0; i <= n; ++i) {</pre>
   if (!z && q[i] == -1) continue;
                                                                    ? (RE(a) < 0 ? 3 : 1) : (IM(a) < 0 ? 0 : 2));
   if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
                                                                  int qb = (IM(b) == 0
                                                                     ? (RE(b) < 0 ? 3 : 1) : (IM(b) < 0 ? 0 : 2));
  if (d[x][s] > -eps) return true;
                                                                  if (qa != qb)
                                                                   return sgn(qa - qb);
  for (int i = 0; i < m; ++i) {
                                                                  return sgn(cross(b, a));
   if (d[i][s] < eps) continue;
if (r == -1 || \</pre>
                                                                 template <typename V> llf area(const V & pt) {
    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
                                                                  11d ret = 0;
                                                                  for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
  if (r == -1) return false;
                                                                   ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
                                                                  return ret / 2.0;
  pivot(r, s);
                                                                 Point rot90(Point p) { return Point{-IM(p), RE(p)}; }
                                                                 Pointf project(Pointf p, Pointf q) { // p onto q
VD solve(const VVD &a, const VD &b, const VD &c) {
m = b.size(), n = c.size();
                                                                  return dot(p, q) * q / dot(q, q);
 d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)
                                                                 6.2 Segment & Line Intersection
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
 p.resize(m), q.resize(n + 1);
                                                                 struct Segment {
 for (int i = 0; i < m; ++i)
                                                                  Point st, dir; // represent st + t*dir for 0 <= t <= 1
 p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
                                                                  Segment(Point s, Point e) : st(s), dir(e - s) {}
 for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i]; q[n] = -1, d[m + 1][n] = 1;
                                                                  static bool valid(lld p, lld q)
                                                                   // is there t s.t. 0 <= t <= 1 && qt == p ?
 int r = 0;
                                                                   if (q < 0) q = -q, p = -p;
 for (int i = 1; i < m; ++i)
if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
                                                                   return 0 <= p && p <= q;
 if (d[r][n + 1] < -eps) {</pre>
                                                                 bool isInter(Segment A, Point P) {
  pivot(r, n)
  if (!phase(1) || d[m + 1][n + 1] < -eps)
                                                                  if (A.dir == Point(0)) return P == A.st;
                                                                  return cross(P - A.st, A.dir) == 0 &&
Segment::valid(dot(P - A.st, A.dir), norm(A.dir));
   return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
   int s = min_element(d[i].begin(), d[i].end() - 1)
         d[i].begin();
                                                                 template <typename U, typename V>
                                                                 bool isInter(U A, V B) {
  if (cross(A.dir, B.dir) == 0)
   pivot(i, s);
                                                                   return // handle parallel yourself
 if (!phase(0)) return VD(n, inf);
                                                                    isInter(A, B.st) || isInter(A, B.st+B.dir) ||
 VD x(n);
                                                                    isInter(B, A.st) || isInter(B, A.st+A.dir);
 for (int i = 0; i < m; ++i)
                                                                  Point D = B.st - A.st;
  if (p[i] < n) \times [p[i]] = d[i][n + 1];
                                                                  11d C = cross(A.dir, B.dir);
                                                                  return U::valid(cross(D, A.dir), C) &&
 return x;
}}
                                                                    V::valid(cross(D, B.dir), C);
5.24 Simplex Construction
                                                                 struct Line {
                                                                  Point st, ed, dir;
Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that for all 1 \leq j \leq m,
                                                                  Line (Point s, Point e)
\sum_{1 \le i \le n} A_{ji} x_i \le b_j and x_i \ge 0 for all 1 \le i \le n.
                                                                   : st(s), ed(e), dir(e - s) {}
  1. In case of minimization, let c_i' = -c_i
                                                                 Pointf intersect(const Line &A, const Line &B) {
                                                                  11f t = cross(B.st - A.st, B.dir) /
  2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
                                                                   llf(cross(A.dir, B.dir));
                                                                  return toPointf(A.st) +
  3. \sum_{1 \le i \le n} A_{ji} x_i = b_j
                                                                   Pointf(t) * toPointf(A.dir);
        • \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
        • \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j
                                                                 6.3 2D Convex Hull
                                                                 template<tvpename PT>
  4. If x_i has no lower bound, replace x_i with x_i - x_i'
                                                                 vector<PT> buildConvexHull(vector<PT> d) {
                                                                  sort(ALL(d), [](const PT& a, const PT& b){
```

return tie(a.x, a.y) < tie(b.x, b.y);});</pre>

vector<PT> s(SZ(d)<<1);</pre>

int o = 0;

6 Geometry

6.1 Basic Geometry

multiset<P, cmp_y> s;

```
for(auto p: d) {
                                                                 void solve(P a[], int n) {
                                                                  sort(a, a + n, [](const P& p, const P& q) {
  while(o>=2 && cross(p-s[o-2], s[o-1]-s[o-2])<=0)
   0--
                                                                   return tie(p.x, p.y) < tie(q.x, q.y);</pre>
  s[o++] = p;
                                                                  11f d = INF; int pt = 0;
 for(int i=SZ(d)-2, t = o+1;i>=0;i--){
                                                                  for (int i = 0; i < n; ++i) {
  while(o = t\&cross(d[i] - s[o-2], s[o-1] - s[o-2]) <= 0)
                                                                   while (pt < i \text{ and } a[i].x - a[pt].x >= d)
                                                                   s.erase(s.find(a[pt++]));
  s[o++] = d[i];
                                                                   auto it = s.lower_bound(P(a[i].x, a[i].y - d));
                                                                   while (it != s.end() and it->y - a[i].y < d)
                                                                    d = min(d, dis(*(it++), a[i]));
 s.resize(o-1);
 return s;
                                                                   s.insert(a[i]);
                                                                }
6.4
      3D Convex Hull
                                                                       kD Closest Pair (3D ver.)
                                                                6.7
// return the faces with pt indexes
int flag[MXN][MXN];
                                                                llf solve(vector<P> v) {
                                                                 shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,</pre>
struct Point{
 ld x,y,z;
 Point operator * (const ld &b) const {
                                                                   unordered_map<lld, int>>> m;
                                                                  llf d = dis(v[0], v[1]);
  return (Point) {x*b,y*b,z*b};}
 Point operator * (const Point &b) const {
                                                                  auto Idx = [&d] (11f x) -> 11d {
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
                                                                   return round(x * 2 / d) + 0.1; };
                                                                  auto rebuild_m = [&m, &v, &Idx](int k) {
                                                                   m.clear();
Point ver(Point a, Point b, Point c) {
  return (b - a) * (c - a);}
                                                                   for (int i = 0; i < k; ++i)
                                                                    m[Idx(v[i].x)][Idx(v[i].y)]
vector<Face> convex_hull_3D(const vector<Point> pt) {
                                                                     [Idx(v[i].z)] = i;
                                                                  }; rebuild_m(2);
 int n = SZ(pt), ftop = 0;
 REP(i,n) REP(j,n) flag[i][j] = 0;
                                                                  for (size_t i = 2; i < v.size(); ++i) {</pre>
                                                                   const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
  kz = Idx(v[i].z); bool found = false;
 vector<Face> now;
 now.emplace_back(0,1,2);
 now.emplace_back(2,1,0);
                                                                   for (int dx = -2; dx <= 2; ++dx) {
                                                                    const 11d nx = dx + kx;
 for (int i=3; i<n; i++){
  ftop++; vector<Face> next;
                                                                    if (m.find(nx) == m.end()) continue;
  REP(j, SZ(now)) {
                                                                    auto& mm = m[nx];
   Face& f=now[j]; int ff = 0;
ld d=(pt[i]-pt[f.a]).dot(
                                                                    for (int dy = -2; dy <= 2; ++dy) {
                                                                     const 11d ny = dy + ky;
     ver(pt[f.a], pt[f.b], pt[f.c]));
                                                                     if (mm.find(ny) == mm.end()) continue;
   if (d <= 0) next.push_back(f);</pre>
                                                                     auto& mmm = mm[ny];
   if (d > 0) ff=ftop;
                                                                     for (int dz = -2; dz <= 2; ++dz) {
   else if (d < 0) ff=-ftop;
                                                                      const 11d nz = dz + kz;
   flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
                                                                      if (mmm.find(nz) == mmm.end()) continue;
                                                                      const int p = mmm[nz];
                                                                      if (dis(v[p], v[i]) < d) {</pre>
  REP(j, SZ(now)) {
   Face& f=now[j]
                                                                       d = dis(v[p], v[i]);
   if (flag[f.a][f.b] > 0 &&
  flag[f.a][f.b] != flag[f.b][f.a])
                                                                       found = true;
    next.emplace_back(f.a,f.b,i);
   if (flag[f.b][f.c] > 0 &&
     flag[f.b][f.c] != flag[f.c][f.b])
    next.emplace_back(f.b,f.c,i);
                                                                   if (found) rebuild_m(i + 1);
   if (flag[f.c][f.a] > 0 &&
                                                                   else m[kx][ky][kz] = i;
     flag[f.c][f.a] != flag[f.a][f.c])
    next.emplace_back(f.c,f.a,i);
                                                                  return d;
  now=next;
                                                                      Simulated Annealing
 return now;
                                                                11f anneal() {
                                                                 mt19937 rnd_engine( seed );
                                                                  uniform_real_distribution< llf > rnd( 0, 1 );
6.5
      2D Farthest Pair
                                                                  const llf dT = 0.001;
// stk is from convex hull
                                                                  // Argument p
                                                                  llf S_cur = calc( p ), S_best = S_cur;
for ( llf T = 2000 ; T > EPS ; T -= dT ) {
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {</pre>
                                                                   // Modify p to p_prime
 while(abs(cross(stk[i+1]-stk[i],
                                                                   const llf S_prime = calc( p_prime );
                                                                   const llf delta_c = S_prime - S_cur;
llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
   stk[(pos+1)%n]-stk[i])) >
   abs(cross(stk[i+1]-stk[i],
   stk[pos]-stk[i]))) pos = (pos+1)%n;
                                                                   if ( rnd( rnd_engine ) <= prob )</pre>
ans = max({ans, dis(stk[i], stk[pos]),
dis(stk[i+1], stk[pos])});
                                                                    S_cur = S_prime, p = p_prime;
                                                                   if ( S_prime < S_best ) // find min</pre>
                                                                    S_best = S_prime, p_best = p_prime;
6.6 2D Closest Pair
                                                                  return S_best;
struct cmp_y {
bool operator()(const P& p, const P& q) const {
                                                                      Half Plane Intersection
  return p.y < q.y;</pre>
                                                                // cross(pt-line.st, line.dir)<=0 <-> pt in half plane
                                                                bool operator<(const Line &lhs, const Line &rhs) {</pre>
```

if (int cmp = argCmp(lhs.dir, rhs.dir))

pdd ft = p1 + project(c-p1, p2-p1), vec = p2-p1;

```
return cmp == -1;
                                                               llf dis = abs(c - ft);
  return ori(lhs.st, lhs.ed, rhs.st) < 0;
                                                               if (abs(dis - r) < eps) return {ft};</pre>
                                                               if (dis > r) return {};
                                                               vec = vec * sqrt(r * r - dis * dis) / abs(vec);
                                                               return {ft + vec, ft - vec};
// intersect function is in "Segment Intersect"
                                                              }
11f HPI(vector<Line> &lines) {
  sort(lines.begin(), lines.end());
                                                                     Intersection of Polygon and Circle
  deque<Line> que;
  deque<Pointf> pt;
                                                              // Divides into multiple triangle, and sum up
  que.push_back(lines[0]);
                                                              // test by HDU2892
  for (int i = 1; i < (int)lines.size(); i++) {
                                                              const double PI=acos(-1);
    if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
                                                              double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
     continue
#define POP(L, R) \
                                                               if(abs(pb)<eps) return 0;</pre>
    while (pt.size() > 0 \
                                                               double S, h, theta;
      && ori(L.st, L.ed, pt.back()) < 0) \
                                                               double a=abs(pb),b=abs(pa),c=abs(pb-pa);
    pt.pop_back(), que.pop_back(); \
while (pt.size() > 0 \
                                                               double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
                                                               double cosC = dot(pa,pb) / a / b, C = acos(cosC);
      && ori(R.st, R.ed, pt.front()) < 0) \
                                                               if(a > r){
      pt.pop_front(), que.pop_front();
                                                                S = (C/2)*r*r;
    POP(lines[i], lines[i]);
                                                                h = a*b*sin(C)/c;
    pt.push_back(intersect(que.back(), lines[i]));
                                                                if (h < r && B < PI/2)</pre>
    que.push_back(lines[i]);
                                                                 S = (acos(h/r)*r*r - h*sqrt(r*r-h*h));
 POP(que.front(), que.back())
                                                               else if(b > r){
 if (que.size() <= 1 ||</pre>
                                                                theta = PI - B - asin(sin(B)/r*a);
    argCmp(que.front().dir, que.back().dir) == 0)
                                                                S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
 pt.push_back(intersect(que.front(), que.back()));
                                                               else S = .5*sin(C)*a*b;
  return area(pt);
                                                               return S;
                                                              double area_poly_circle(const vector<pdd> &poly,
6.10 Minkowski Sum
                                                                const pdd &0,const double r){
                                                               double S=0; int N=poly.size();
for(int i=0;i<N;++i)</pre>
vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
hull(A), hull(B);
vector<pll> C(1, A[0] + B[0]), s1, s2;
                                                                S += _area(poly[i]-0, poly[(i+1)%N]-0, r)
for(int i = 0; i < SZ(A); ++i)
s1.pb(A[(i + 1) % SZ(A)] - A[i]);
                                                                  * ori(0, poly[i], poly[(i+1)%N]);
                                                               return fabs(S);
for(int i = 0; i < SZ(B); i++)
 s2.pb(B[(i + 1) % SZ(B)] - B[i]);
                                                                     Tangent line of Two Circle
                                                              6.14
 for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
 if (p2 >= SZ(B)
                                                              vector<Line> go(const Cir &c1, const Cir &c2,
    || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
                                                                int sign1) {
   C.pb(C.back() + s1[p1++]);
                                                               // sign1 = 1 for outer tang, -1 for inter tang
                                                               vector<Line> ret;
  C.pb(C.back() + s2[p2++]);
                                                               if (norm(c1.o - c2.o) < eps)
 return hull(C), C;
                                                                return ret
                                                               11f d = abs(c1.o - c2.o);
                                                               Pointf v = (c2.o - c1.o) / d;
6.11 Circle Class
                                                               11f c = (c1.r - sign1 * c2.r) / d;
struct Circle { Pointf o; llf r; };
                                                               if(c*c>1)
                                                                return ret;
vector<llf> intersectAngle(Circle A, Circle B) {
                                                               11f h = sqrt(max(0.0, 1.0 - c * c));
Pointf dir = B.o - A.o; llf d2 = norm(dir);
                                                               for (int sign2: {1, -1}) {
                                                                Pointf n = c * v + sign2 * h * rot90(v);
if (norm(A.r - B.r) >= d2)
 if (A.r < B.r) return {-PI, PI}; // special</pre>
                                                                Pointf p1 = c1.o + n * c1.r;
  else return {};
                                                                Pointf p2 = c2.0 + n * (c2.r * sign1);
 if (norm(A.r + B.r) <= d2) return {};</pre>
                                                                if (norm(p2 - p1) < eps)
llf dis = abs(dir), theta = arg(dir);
llf phi = acos((A.r * A.r + d2 - B.r * B.r) /
                                                                 p2 = p1 + rot90(c2.o - c1.o);
                                                                ret.push_back({p1, p2});
   (2 * A.r * dis));
11f L = theta - phi, R = theta + phi;
                                                               return ret;
while (L < -PI) L += PI * 2;
while (R > PI) R -= PI * 2;
                                                              6.15
                                                                     Minimum Covering Circle
return { L, R };
                                                              template<typename P>
                                                              Circle getCircum(const P &a, const P &b, const P &c){
vector<Pointf> intersectPoint(Circle a, Circle b) {
                                                               Real a1 = a.x-b.x, b1 = a.y-b.y;
11f d = abs(a.o - b.o);
                                                               Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
if (d >= b.r+a.r || d <= abs(b.r-a.r)) return {};</pre>
                                                               Real a2 = a.x-c.x, b2 = a.y-c.y;
Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
11f dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
Pointf dir = (a.o - b.o); dir /= d;
                                                               Circle cc;
Pointf pcrs = dir*d1 + b.o;
                                                               cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
dt=sqrt(max(0.0L, b.r*b.r-d1*d1)), dir = rot90(dir);
                                                               cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2)
return {pcrs + dir*dt, pcrs - dir*dt};
                                                               cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
                                                               return cc;
6.12 Intersection of line and Circle
vector<pdd> line_interCircle(const pdd &p1,
                                                              template<typename P>
  const pdd &p2, const pdd &c, const double r) {
                                                              Circle MinCircleCover(const vector<P>& pts){
```

random_shuffle(pts.begin(), pts.end());

```
Circle c = { pts[0], 0 };
                                                                   } else {
for(int i=0;i<(int)pts.size();i++){</pre>
                                                                    nearest(r->R, x, y, mID, md2);
 if (dist(pts[i], c.o) <= c.r) continue;</pre>
                                                                    nearest(r->L, x, y, mID, md2);
 c = { pts[i], 0 };
for (int j = 0; j < i; j++) {</pre>
   if(dist(pts[j], c.o) <= c.r) continue;</pre>
                                                                  int query(int x, int y) {
   c.o = (pts[i] + pts[j]) / 2;
                                                                   int id = 1029384756;
                                                                   LL d2 = 102938475612345678LL;
   c.r = dist(pts[i], c.o);
   for (int k = 0; k < j; k++) {
                                                                   nearest(root, x, y, id, d2);
    if (dist(pts[k], c.o) <= c.r) continue;</pre>
                                                                   return id:
    c = getCircum(pts[i], pts[j], pts[k]);
                                                                } tree;
 }
                                                                 6.17 Rotating Sweep Line
return c;
                                                                 void rotatingSweepLine(pair<int, int> a[], int n) {
                                                                  vector<pair<int, int>> 1;
                                                                  1.reserve(n * (n - 1) / 2);
      KDTree (Nearest Point)
                                                                  for (int i = 0; i < n; ++i)
                                                                   for (int j = i + 1; j < n; ++j)
const int MXN = 100005;
struct KDTree {
                                                                    1.emplace_back(i, j);
                                                                  sort(1.begin(), 1.end(), [&a](auto &u, auto &v){
struct Node {
  int x,y,x1,y1,x2,y2;
                                                                   11d udx = a[u.first].first - a[u.second].first;
                                                                   lld udy = a[u.first].second - a[u.second].second;
lld vdx = a[v.first].first - a[v.second].first;
 int id,f;
Node *L, *R;
                                                                   11d vdy = a[v.first].second - a[v.second].second;
 } tree[MXN], *root;
                                                                   if (udx == 0 \text{ or } vdx == 0) \text{ return not } udx == 0;
 int n:
                                                                   int s = sgn(udx * vdx);
LL dis2(int x1, int y1, int x2, int y2) {
 LL dx = x1-x2, dy = y1-y2;
                                                                   return udy * vdx * s < vdy * udx * s;</pre>
 return dx*dx+dy*dy;
                                                                  });
                                                                  vector<int> idx(n), p(n);
                                                                  iota(idx.begin(), idx.end(), 0);
static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
                                                                  sort(idx.begin(), idx.end(), [&a](int i, int j){
 static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
                                                                   return a[i] < a[j]; });</pre>
 void init(vector<pair<int,int>> ip) {
                                                                  for (int i = 0; i < n; ++i) p[idx[i]] = i;</pre>
 n = ip.size();
 for (int i=0; i<n; i++) {</pre>
                                                                  for (auto [i, j]: 1) {
   tree[i].id = i;
tree[i].x = ip[i].first;
                                                                   // do here
                                                                   swap(p[i], p[j]);
   tree[i].y = ip[i].second;
                                                                   idx[p[i]] = i, idx[p[j]] = j;
  root = build_tree(0, n-1, 0);
                                                                 6.18
                                                                        Circle Cover
Node* build_tree(int L, int R, int d) {
                                                                 const int N = 1021;
 if (L>R) return nullptr
  int M = (L+R)/2; tree[M].f = d%2;
                                                                 struct CircleCover {
  nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
                                                                  int C;
  tree[M].x1 = tree[M].x2 = tree[M].x;
                                                                  Cir c[N]
 tree[M].y1 = tree[M].y2 = tree[M].y;
                                                                  bool g[N][N], overlap[N][N];
  tree[M].L = build_tree(L, M-1, d+1);
                                                                  // Area[i] : area covered by at least i circles
double Area[ N ];
  if (tree[M].L) {
  tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
                                                                  void init(int _C){ C = _C;}
   tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
                                                                  struct Teve {
                                                                   pdd p; double ang; int add;
Teve() {}
   tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
                                                                   Teve(pdd _a, double _b, int _c):p(_a), ang(_b), add(
  tree[M].R = build_tree(M+1, R, d+1);
                                                                      c){}
  if (tree[M].R) {
                                                                   bool operator<(const Teve &a)const
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
                                                                   {return ang < a.ang;}
   tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
                                                                  }eve[N * 2];
                                                                   // strict: x = 0, otherwise x = -1
   tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
                                                                  bool disjuct(Cir &a, Cir &b, int x)
                                                                  {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
bool contain(Cir &a, Cir &b, int x)
  }
  return tree+M;
                                                                  {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
 int touch(Node* r, int x, int y, LL d2){
                                                                  bool contain(int i, int j) {
 LL dis = sqrt(d2)+1;
                                                                   /* c[j] is non-strictly in c[i]. */
 if (x<r->x1-dis || x>r->x2+dis ||
                                                                   return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c[i].R) \mid c[i].R - c[i].R)
    y<r->y1-dis || y>r->y2+dis)
                                                                      [j].R) == 0 \&\& i < j)) \&\& contain(c[i], c[j], -1);
   return 0;
                                                                  void solve(){
 return 1;
                                                                   fill_n(Area, C + 2, 0);
                                                                   for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)</pre>
void nearest(Node* r,int x,int y,int &mID,LL &md2) {
 if (!r || !touch(r, x, y, md2)) return;
 LL d2 = dis2(r->x, r->y, x, y);
                                                                     overlap[i][j] = contain(i, j);
                                                                   for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j</pre>
  if (d2 < md2 \mid | (d2 == md2 \&\& mID < r->id)) {
  mID = r->id;
                                                                      g[i][j] = !(overlap[i][j] || overlap[j][i] ||
   md2 = d2;
                                                                        disjuct(c[i], c[j], -1));
  // search order depends on split dim
                                                                   for(int i = 0; i < C; ++i){
  if ((r->f == 0 && x < r->x) ||
                                                                    int E = 0, cnt = 1;
    (r->f == 1 && y < r->y))
                                                                    for(int j = 0; j < C; ++j)
   nearest(r->L, x, y, mID, md2);
                                                                     if(j != i && overlap[j][i])
   nearest(r->R, x, y, mID, md2);
                                                                      ++cnt;
```

```
for(int j = 0; j < C; ++j)
    if(i != j && g[i][j]) {
     pdd aa, bb;
     CCinter(c[i], c[j], aa, bb);

llf A = atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);

llf B = atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);
     eve[E++] = Teve(bb,B,1), eve[E++]=Teve(aa,A,-1);
     if(B > A) ++cnt;
   if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
   else{
    sort(eve, eve + E);
    eve[E] = eve[0];
    for(int j = 0; j < E; ++j){
     cnt += eve[j].add;
     Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
     double theta = eve[j + 1].ang - eve[j].ang;
     if (theta < 0) theta += 2. * pi;</pre>
     Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
};
7
     Stringology
```

7.1 Hash

```
class Hash {
  private:
    static constexpr int P = 127, Q = 1051762951;
    vector<int> h, p;
  public:
    void init(const string &s) {
      h.assign(s.size()+1, 0); p.resize(s.size()+1);
      for (size_t i = 0; i < s.size(); ++i)
         h[i + 1] = add(mul(h[i], P), s[i]);
      generate(p.begin(), p.end(), [x=1,y=1,this]()
         mutable{y=x;x=mul(x,P);return y;});
    }
    int query(int l, int r) { // 1-base (l, r]
      return sub(h[r], mul(h[1], p[r-1]));}
};</pre>
```

7.2 Suffix Array

```
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
int x[maxn], p[maxn], q[maxn * 2];
// sa[i]: sa[i]-th suffix is the \
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix \
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *sa, int *c, int n, int z) {
 memset(sa, 0, sizeof(int) * n);
 memcpy(x, c, sizeof(int) * z);
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
 if (sa[i] && !t[sa[i] - 1])
   sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
 memcpy(x, c, sizeof(int) * z);
 for (int i = n - 1; i >= 0; --i)
  if (sa[i] && t[sa[i] - 1])
   sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q,
bool *t, int *c, int n, int z) {
bool uniq = t[n - 1] = true;
 int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
 memset(c, 0, sizeof(int) * z);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
 for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
 if (uniq) {
  for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
  return;
 for (int i = n - 2; i >= 0; --i)
 t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
```

```
pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i)</pre>
  if (t[i] && !t[i - 1])
    sa[--x[s[i]]] = p[q[i] = nn++] = i;
 induce(sa, c, s, t, n, z);

for (int i = 0; i < n; ++i) {

  if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
  bool neq = last < 0 || \
   memcmp(s + sa[i], s + last,
(p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
  ns[q[last = sa[i]]] = nmxz += neq;
 sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
 pre(sa, c, n, z);
 for (int i = nn - 1; i >= 0; --i)
  sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
 induce(sa, c, s, t, n, z);
void build(const string &s) {
 for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
 _s[(int)s.size()] = 0; // s shouldn't contain 0
 sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];
for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
 int ind = 0; hi[0] = 0;
 for (int i = 0; i < (int)s.size(); ++i) {</pre>
  if (!rev[i]) {
   ind = 0;
    continue;
  while (i + ind < (int)s.size() && \</pre>
   s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  hi[rev[i]] = ind ? ind-- : 0;
}}
      Suffix Automaton
```

```
struct SuffixAutomaton {
 struct node +
  int ch[K], len, fail, cnt, indeg;
  node(int L = 0) : ch{}, len(L), fail(0), cnt(0),
    indeg(0) {}
 } st[N];
 int root, last, tot;
 void extend(int c) {
  int cur = ++tot;
  st[cur] = node(st[last].len + 1);
  while (last && !st[last].ch[c]) {
    st[last].ch[c] = cur;
    last = st[last].fail;
  if (!last) {
    st[cur].fail = root;
    int q = st[last].ch[c];
    if (st[q].len == st[last].len + 1) {
      st[cur].fail = q;
    } else {
      int clone = ++tot;
      st[clone] = st[q];
      st[clone].len = st[last].len + 1;
      st[st[cur].fail = st[q].fail = clone].cnt = 0;
      while (last && st[last].ch[c] == q) {
        st[last].ch[c] = clone;
        last = st[last].fail;
      }
   }
 st[last = cur].cnt += 1;
 void init(const char* s) {
 root = last = tot = 1;
  st[root] = node(0);
  for (char c; c = *s; ++s) extend(c - 'a');
 int q[N];
 void dp()
  for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
  int head = 0, tail = 0;
  for (int i = 1; i <= tot; i++)</pre>
    if (st[i].indeg == 0) q[tail++] = i;
```

```
while (head != tail) {
                                                                 int k = 0;
                                                                 while (k < n \&\& s[i+k] == s[j+k]) k++;
    int now = q[head++];
    if (int f = st[now].fail) {
                                                                 if (s[i+k] <= s[j+k]) j += k+1;
      st[f].cnt += st[now].cnt;
                                                                 else i += k+1;
      if (--st[f].indeg == 0) q[tail++] = f;
                                                                 if (i == j) j++;
 }
                                                                int ans = i < n ? i : j;</pre>
}
                                                                return s.substr(ans, n);
int run(const char* s) {
  int now = root;
                                                               7.8 BWT
 for (char c; c = *s; ++s) {
   if (!st[now].ch[c -= 'a']) return 0;
                                                               struct BurrowsWheeler{
   now = st[now].ch[c];
                                                               #define SIGMA 26
                                                               #define BASE 'a'
  return st[now].cnt;
                                                                vector<int> v[ SIGMA ];
                                                                void BWT(char* ori, char* res){
}
} SAM;
                                                                 // make ori -> ori + ori
                                                                 // then build suffix array
7.4
     KMP
                                                                void iBWT(char* ori, char* res){
vector<int> kmp(const string &s) {
                                                                 for( int i = 0 ; i < SIGMA ; i ++ )</pre>
vector<int> f(s.size(), θ);
                                                                  v[ i ].clear();
 /* f[i] = length of the longest prefix
                                                                 int len = strlen( ori );
for( int i = 0 ; i < len ; i ++ )</pre>
   (excluding s[0:i]) such that it coincides
   with the suffix of s[0:i] of the same length */
                                                                  v[ ori[i] - BASE ].push_back( i );
 /* i + 1 - f[i] is the length of the
                                                                 vector<int> a:
   smallest recurring period of s[0:i] */
                                                                 for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
 int k = 0:
                                                                  for( auto j : v[ i ] ){
 for (int i = 1; i < (int)s.size(); ++i) {</pre>
                                                                   a.push_back( j );
ori[ ptr ++ ] = BASE + i;
 while (k > 0 \& s[i] != s[k]) k = f[k - 1];
 if (s[i] == s[k]) ++k;
 f[i] = k;
                                                                 for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
return f;
                                                                  ptr = a[ ptr ];
vector<int> search(const string &s, const string &t) {
                                                                 res[ len ] = 0;
 // return 0-indexed occurrence of t in s
vector<int> f = kmp(t), r;
for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
                                                               } bwt;
 while(k > 0 \&\& (k==(int)t.size() \mid | s[i]!=t[k]))
                                                               7.9 Palindromic Tree
  k = f[k - 1];
                                                               struct palindromic_tree{
 if (s[i] == t[k]) ++k;
                                                                struct node{
  if (k == (int)t.size()) r.push_back(i-t.size()+1);
                                                                 int next[26],f,len;
                                                                 int cnt,num,st,ed;
 return res;
                                                                 node(int l=0):f(0),len(1),cnt(0),num(0) {
                                                                  memset(next, 0, sizeof(next)); }
7.5 Z value
                                                                vector<node> st:
vector<int> Zalgo(const string &s) {
                                                                vector<char> s;
vector<int> z(s.size(), s.size());
                                                                int last,n;
for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
                                                                void init(){
 int j = clamp(r - i, 0, z[i - 1]);
for (; i + j < z[0] and s[i + j] == s[j]; ++j);
                                                                 st.clear();s.clear();last=1; n=0;
                                                                 st.push_back(0);st.push_back(-1);
 if (i + (z[i] = j) > r) r = i + z[1 = i];
                                                                 st[0].f=1;s.push_back(-1); }
                                                                int getFail(int x){
return z;
                                                                 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
                                                                 return x;}
                                                                void add(int c){
7.6 Manacher
                                                                 s.push_back(c-='a'); ++n;
int z[maxn];
                                                                 int cur=getFail(last);
int manacher(const string& s) {
  string t = ".";
                                                                 if(!st[cur].next[c]){
                                                                  int now=st.size();
for(char c: s) t += c, t += '.';
                                                                  st.push_back(st[cur].len+2);
int 1 = 0, r = 0, ans = 0;
                                                                  st[now].f=st[getFail(st[cur].f)].next[c];
for (int i = 1; i < t.length(); ++i) {</pre>
                                                                  st[cur].next[c]=now;
 z[i] = (r > i'] min(z[2 * 1' - i], r - i) : 1);
while (i - z[i] >= 0 && i + z[i] < t.length()) {
                                                                  st[now].num=st[st[now].f].num+1;
  if(t[i - z[i]] == t[i + z[i]]) ++z[i];
                                                                 last=st[cur].next[c];
   else break:
                                                                 ++st[last].cnt;}
                                                                void dpcnt() {
 if (i + z[i] > r) r = i + z[i], l = i;
                                                                 for (int i=st.size()-1; i >= 0; i--)
                                                                  st[st[i].f].cnt += st[i].cnt;
for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
return ans:
                                                                int size(){ return st.size()-2;}
                                                               } pt;
                                                               int main() {
      Lexico Smallest Rotation
                                                                string s; cin >> s; pt.init();
                                                                for (int i=0; i<SZ(s); i++)
string mcp(string s){
int n = s.length();
                                                                 int prvsz = pt.size(); pt.add(s[i]);
                                                                 if (prvsz != pt.size()) -
s += s;
int i=0, j=1;
                                                                  int r = i, l = r - pt.st[pt.last].len + 1;
while (i<n && j<n){</pre>
                                                                  // pal @ [1,r]: s.substr(1, r-1+1)
```

8 Misc

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\dots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \le k \le n$.

8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W\subseteq X, |W|\leq |N_G(W)|$

8.1.7 Euler's planar graph formula

$$V - E + F = C + 1$$
, $E \le 3V - 6$ (?)

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Lucas's theorem

 $\binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p}, \text{ where } m=m_k p^k+m_{k-1} p^{k-1}+\cdots+m_1 p+m_0,$ and $n=n_k p^k+n_{k-1} p^{k-1}+\cdots+n_1 p+n_0.$

8.1.10 Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2),$ find maximum $S\in I_1\cap I_2.$ For each iteration, build the directed graph and find a shortest path from s to t.

- $s \rightarrow x : S \sqcup \{x\} \in I_1$
- $x \rightarrow t : S \sqcup \{x\} \in I_2$
- $y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1$ (y is in the unique circuit of $S \sqcup \{x\}$)
- $x \to y: S \setminus \{y\} \sqcup \{x\} \in I_2$ (y is in the unique circuit of $S \sqcup \{x\}$)

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|.$ In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

8.2 DP-opt Condition

8.2.1 totally monotone (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

8.2.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

8.3 Convex 1D/1D DP

```
struct segment {
int i, 1, r;
segment() {}
segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 11d f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
dp[0] = 0;
 deque<segment> dq; dq.push_back(segment(0, 1, n));
 for (int i = 1; i <= n; ++i) {
 dp[i] = f(dq.front().i, i);
 while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
 dq.front().l = i + 1;
 segment seg = segment(i, i + 1, n);
 while (dq.size() &&
  f(i, dq.back().1) < f(dq.back().i, dq.back().1))
   dq.pop_back();
 if (dq.size()) {
  int d = 1 << 20, c = dq.back().1;</pre>
   while (d \gg 1) if (c + d \ll d, back().r)
   if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.1 = c + 1;
 if (seg.1 <= n) dq.push_back(seg);</pre>
```

8.4 ConvexHull Optimization

```
struct Line {
 mutable int64_t a, b, p;
 bool operator<(const Line &rhs) const { return a < rhs
    .a; }
 bool operator<(int64_t x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>> {
 static const int64_t kInf = 1e18;
 bool Isect(iterator x, iterator y)
  auto Div = [](int64_t a, int64_t b) {
  return a / b - ((a ^ b) < 0 && a % b); }
if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(int64_t a, int64_t b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
  if (x != begin() \&\& Isect(--x, y)) Isect(x, y = erase)
    (y));
  while ((y = x) != begin() && (--x)->p >= y->p) Isect(
    x, erase(y));
 int64_t Query(int64_t x) {
  auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
};
```

8.5 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

8.6 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
```

```
void tarjan(int u){
                                                                 for(int i=0;i<m;i++){</pre>
                                                                 scanf("%d%d",&a,&b)
dfn[u]=low[u]=++dfs_idx;
for(int i=0;i<(int)init_g[u].size();i++){</pre>
                                                                  init_g[a].push_back(b);
  int v=init_g[u][i];
                                                                  init_g[b].push_back(a);
  if(v==par[u]) continue;
  if(!dfn[v]){
                                                                par[1]=-1;
   par[v]=u;
                                                                 tarjan(1);
                                                                dfs(1,-1);
printf("%d\n",max(dp[1][0],dp[1][1]));
   tarjan(v);
   low[u]=min(low[u],low[v]);
   if(dfn[u]<low[v]){</pre>
                                                                 return 0;
    g[u].push_back(v);
    g[v].push_back(u);
                                                               8.7
                                                                     Tree Knapsack
  }else{
                                                               int dp[N][K]; PII obj[N];
   low[u]=min(low[u],dfn[v]);
                                                               vector<int> G[N];
   if(dfn[v]<dfn[u]){</pre>
                                                               void dfs(int u, int mx){
    int temp_v=u;
                                                                 for(int s: G[u]) {
    bcc_id++;
                                                                  if(mx < obj[s].first) continue;</pre>
    while(temp_v!=v){
                                                                  for(int i=0;i<=mx-obj[s].FF;i++)</pre>
     g[bcc_id+n].push_back(temp_v);
                                                                   dp[s][i] = dp[u][i];
     g[temp_v].push_back(bcc_id+n);
                                                                  dfs(s, mx - obj[s].first);
     temp_v=par[temp_v];
                                                                  for(int i=obj[s].FF;i<=mx;i++)</pre>
                                                                   dp[u][i] = max(dp[u][i],
                                                                    dp[s][i - obj[s].FF] + obj[s].SS);
    g[bcc_id+n].push_back(v);
    g[v].push_back(bcc_id+n);
                                                                }
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
                                                               }
                                                                8.8 N Queens Problem
                                                               vector< int > solve( int n ) {
                                                                 // no solution when n=2, 3
int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
                                                                 vector< int > ret;
                                                                if ( n % 6 == 2 ) {
  for ( int i = 2 ; i <= n ; i += 2 )
    ret.push_back( i );</pre>
void dfs(int u,int fa){
 if(u<=n){
  for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                  ret.push_back( 3 ); ret.push_back( 1 );
for ( int i = 7 ; i <= n ; i += 2 )
   int v=g[u][i];
   if(v==fa) continue;
                                                                   ret.push_back( i );
   dfs(v,u);
   memset(tp,0x8f,sizeof tp);
                                                                  ret.push_back( 5 );
                                                                 } else if ( n % 6 == 3 ) {
   if(v<=n){
                                                                  for ( int i = 4 ; i <= n ; i += 2 )
    tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
    tp[1]=max(
                                                                   ret.push_back( i );
                                                                 ret.push_back( 2 );
for ( int i = 5 ; i <= n ; i += 2 )
     dp[u][0]+dp[v][0]+1
     dp[u][1]+max(dp[v][0],dp[v][1])
                                                                   ret.push_back( i );
   }else{
                                                                  ret.push_back( 1 ); ret.push_back( 3 );
    tp[0]=dp[u][0]+dp[v][0];
                                                                 } else {
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
                                                                  for ( int i = 2 ; i <= n ; i += 2 )
                                                                   ret.push_back( i );
   dp[u][0]=tp[0],dp[u][1]=tp[1];
                                                                  for ( int i = 1 ; i <= n ; i += 2 )
                                                                   ret.push_back( i );
 }else{
                                                                 }
  for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                 return ret;
  int v=g[u][i];
                                                               }
   if(v==fa) continue;
                                                                     Aliens Optimization
   dfs(v,u);
                                                               long long Alien() {
  min_dp[0][0]=0;
                                                                 long long c = kInf;
  min_dp[1][1]=1;
                                                                 for (int d = 60; d >= 0; --d) {
  min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
                                                                  // cost can be negative, depending on the problem.
  for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                  if (c - (1LL << d) < 0) continue;</pre>
                                                                 long long ck = c - (1LL << d);
pair<long long, int> r = check(ck);
   int v=g[u][i];
   if(v==fa) continue;
   memset(tmp,0x8f,sizeof tmp);
                                                                  if (r.second == k) return r.first - ck * k;
   tmp[0][0]=max(
                                                                 if (r.second < k) c = ck;
    min_dp[0][0]+max(dp[v][0],dp[v][1]),
    min_dp[0][1]+dp[v][0]
                                                                pair<long long, int> r = check(c);
                                                                 return r.first - c * k;
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
   tmp[1][0]=max(
                                                               8.10 Hilbert Curve
   min_dp[1][0]+max(dp[v][0],dp[v][1]),
    min_dp[1][1]+dp[v][0]
                                                               long long hilbert(int n, int x, int y) {
                                                                 long long res = 0;
                                                                 for (int s = n / 2; s; s >>= 1) {
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
                                                                 int rx = (x & s) > 0, ry = (y & s) > 0;
res += s * 111 * s * ((3 * rx) ^ ry);
  memcpy(min_dp,tmp,sizeof tmp);
  dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
                                                                  if (ry == 0) {
  dp[u][0]=min_dp[0][0];
                                                                   if (rx == 1) x = s - 1 - x, y = s - 1 - y;
                                                                   swap(x, y);
int main(){
int m,a,b;
                                                                 return res;
scanf("%d%d",&n,&m);
```

8.11 Binary Search On Fraction

```
struct Q {
11 p, q;
Q go(Q b, 11 d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
 Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
 assert(pred(hi));
 bool dir = 1, L = 1, H = 1;
for (; L || H; dir = !dir) {
  11 len = 0, step = 1;
for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
   if (Q mid = hi.go(lo, len + step);
     mid.p > N || mid.q > N || dir ^ pred(mid))
    t++;
   else len += step;
  swap(lo, hi = hi.go(lo, len));
(dir ? L : H) = !!len;
 return dir ? hi : lo;
```