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1.1 vimrc

```
se is nu ru et tgc sc hls cin cino+=j1 sw=4 sts=4 bs=2
    mouse=a "encoding=utf-8 ls=2
syn on
colo desert
filetype indent on
inoremap {<CR> {<CR>}<ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -std=c++17 -
    DCKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
    Wconversion -fsanitize=address,undefined -g && echo
     success<CR>
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 &&
     echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
```

1.2 Debug Macro

```
#ifdef CKISEKI
#define safe cerr<<__PRETTY_FUNCTION__\
<<" line "<<__LINE__<<" safe\n"</pre>
#define debug(a...) debug_(#a, a)
#define orange(a...) orange_(#a, a)
using std::cerr;
template <typename ...T>
void debug_(const char *s, T ...a) {
  cerr << "\e[1;32m(" << s << ") = (";
  int cnt = sizeof...(T);
  (..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
template <typename Iter>
void orange_(const char *s, Iter L, Iter R) {
  cerr << "\e[1;32m[ " << s << " ] = [ ";

for (int f = 0; L != R; ++L)

cerr << (f++ ? ", " : "") << *L;
  cerr << " ]\e[0m\n";</pre>
#else
#define safe ((void)0)
#define debug(...) ((void)0)
#define orange(...) ((void)0)
#endif
```

1.3 Increase Stack

```
const int size = 256 << 20;</pre>
                                    register long rsp asm("rsp");
                                    char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
                                    __asm__("movq %0, %%rsp\n"::"r"(p));
                                    // main
```

1.4 Pragma Optimization

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8000)
```

1.5 IO Optimization

```
static inline int gc() {
constexpr int B = 1<<20;</pre>
static char buf[B], *p, *q;
if(p == q &&
  (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
  return EOF:
return *p++;
```

2 **Data Structure**

2.1 Dark Magic

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using namespace __gnu_pbds;
// heap tags: paring/binary/binomial/rc_binomial/thin
template<typename T>
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>, \
                  pairing_heap_tag>;
// pbds_heap::point_iterator
// x = pq.push(10); pq.modify(x, 87); a.join(b);
// tree tags: rb_tree_tag/ov_tree_tag/splay_tree_tag
template<typename T>
using ordered_set = tree<T, null_type, less<T>,
  rb_tree_tag, tree_order_statistics_node_update>;
  find_by_order, order_of_key
// hash tables: cc_hash_table/gp_hash_table
```

2.2 Link-Cut Tree

```
template <typename Val, typename SVal> class LCT {
 struct node {
  int pa, ch[2];
  bool rev;
  Val v, prod, rprod;
  SVal sv, sub, vir;
  node() : pa{0}, ch{0, 0}, rev{false}, v{}, prod{},
    rprod{}, sv{}, sub{}, vir{} {};
#define cur o[u]
#define lc cur.ch[0]
#define rc cur.ch[1]
 vector<node> o;
 bool is_root(int u) const {
  return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u;
 bool is_rch(int u) const {
  return o[cur.pa].ch[1] == u && !is_root(u);
 void down(int u) {
  if (not cur.rev) return;
  if (lc) set_rev(lc);
  if (rc) set_rev(rc);
  cur.rev = false;
 void up(int u) {
  cur.prod = o[lc].prod * cur.v * o[rc].prod;
 cur.rprod = o[rc].rprod * cur.v * o[lc].rprod;
  cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
 void set_rev(int u) {
  swap(lc, rc);
  swap(cur.prod, cur.rprod);
  cur.rev ^= 1;
 void rotate(int u) {
  int f=cur.pa, g=o[f].pa, l=is_rch(u);
if (cur.ch[l ^ 1]) o[cur.ch[l ^ 1]].pa = f;
  if (not is_root(f)) o[g].ch[is_rch(f)] = u;
  o[f].ch[l] = cur.ch[l ^ 1];
  cur.ch[l ^ 1] = f;
  cur.pa = g, o[f].pa = u;
  up(f);
```

```
void splay(int u) {
  vector<int> stk = {u};
  while (not is_root(stk.back()))
   stk.push_back(o[stk.back()].pa);
  while (not stk.empty()) {
   down(stk.back());
   stk.pop_back();
  for (int f = cur.pa; not is_root(u); f = cur.pa) {
   if(!is_root(f))rotate(is_rch(u)==is_rch(f)?f:u);
   rotate(u);
 up(u);
 void access(int x) {
  for (int u = x, last = 0; u; u = cur.pa) {
   splay(u);
   cur.vir = cur.vir + o[rc].sub - o[last].sub;
  rc = last; up(last = u);
 splay(x);
 int find_root(int u) {
  int la = 0;
  for (access(u); u; u = lc) down(la = u);
  return la;
 void split(int x, int y) {change_root(x);access(y);}
 void change_root(int u) { access(u); set_rev(u); }
public:
 LCT(int n = 0) : o(n + 1) {}
 int add(const Val &v = {}) {
 o.push_back(v);
 return int(o.size()) - 2;
 int add(Val &&v) {
 o.emplace_back(move(v));
  return int(o.size()) - 2;
 void set_val(int u, const Val &v) {
  splay(++u); cur.v = v; up(u);
 void set_sval(int u, const SVal &v) {
  splay(++u); cur.sv = v; up(u);
 Val query(int x, int y) {
 split(++x, ++y); return o[y].prod;
 SVal subtree(int p, int u) {
 change_root(++p); access(++u);
  return cur.vir + cur.sv;
 bool connected(int u, int v) {
  return find_root(++u) == find_root(++v); }
 void link(int x, int y) {
 change_root(++x); access(++y);
  o[y].vir = o[y].vir + o[x].sub;
 up(o[x].pa = y);
 void cut(int x, int y) {
 split(++x, ++y);
  o[y].ch[0] = o[x].pa = 0; up(y);
#undef cur
#undef lc
#undef rc
2.3 LiChao Segment Tree
```

```
struct L {
 int m, k, id;
L() : id(-1) {}
 L(int a, int b, int c) : m(a), k(b), id(c) {}
 int at(int x) { return m * x + k; }
class LiChao {
private:
 int n; vector<L> nodes;
 static int lc(int x) { return 2 * x + 1; }
 static int rc(int x) { return 2 * x + 2; }
 void insert(int l, int r, int id, L ln) {
  int m = (l + r) >> 1;
```

```
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  if (nodes[id].id == -1) {
  nodes[id] = ln;
  return;
 bool atLeft = nodes[id].at(l) < ln.at(l);</pre>
 if (nodes[id].at(m) < ln.at(m)) {</pre>
  atLeft ^= 1;
  swap(nodes[id], ln);
 if (r - l == 1) return;
 if (atLeft) insert(l, m, lc(id), ln);
 else insert(m, r, rc(id), ln);
int query(int l, int r, int id, int x) {
 int ret = 0, m = (l + r) >> 1;
 if (nodes[id].id != -1)
  ret = nodes[id].at(x);
 if (r - l == 1) return ret;
 if (x < m) return max(ret, query(l, m, lc(id), x));</pre>
  return max(ret, query(m, r, rc(id), x));
public:
LiChao(int n_{-}): n(n_{-}), nodes(n * 4) {}
void insert(L ln) { insert(0, n, 0, ln); }
int query(int x) { return query(0, n, 0, x); }
2.4 Treap
namespace Treap{
#define sz(x)((x)?((x)->size):0)
struct node{
 int size; uint32_t pri;
 node *lc, *rc, *pa;
 node():size(0),pri(rnd()),lc(0),rc(0),pa(0){}
 void pull() {
  size = 1; pa = nullptr;
  if ( lc ) { size += lc->size; lc->pa = this; }
  if ( rc ) { size += rc->size; rc->pa = this; }
 }
};
node* merge( node* L, node* R ) {
 if ( not L or not R ) return L ? L : R;
 if ( L->pri > R->pri ) {
  L->rc = merge( L->rc, R ); L->pull();
  return L;
 } else {
  R->lc = merge( L, R->lc ); R->pull();
  return R;
 }
```

2.5 Linear Basis

L->pull();

R->pull();

 $} // sz(L) == k$

return r;

#undef sz

}

} else {
 R = rt;

```
template <int BITS, typename S = int> struct Basis {
   static constexpr auto MIN = numeric_limits<S>::min();
   array<pair<uint64_t, S>, BITS> b;
   Basis() { b.fill({0, MIN}); }
   void add(uint64_t x, S p) {
      for (int i = BITS-1; i>=0; i--) if ((x >> i) & 1) {
        if (b[i].first == 0) return b[i]={x, p}, void();
        if (b[i].second < p)</pre>
```

void split_by_size(node*rt,int k,node*&L,node*&R) {

split_by_size(rt->rc,k-sz(rt->lc)-1,L->rc,R);

if (not rt) L = R = nullptr;

int getRank(node *o) { // 1-base

for (;o->pa != nullptr; o = o->pa)

int r = sz(o->lc) + 1;

else if(sz(rt->lc) + 1 <= k) {

split_by_size(rt->lc, k, L, R->lc);

if (o->pa->rc == o) r += sz(o->pa->lc) + 1;

```
swap(b[i].first, x), swap(b[i].second, p);
    x ^= b[i].first;
  }
  optional<uint64_t> query_kth(uint64_t v, uint64_t k){
   vector<pair<uint64_t, int>> o;
   for (int i = 0; i < BITS; i++)</pre>
    if (b[i].first) o.emplace_back(b[i].first, i);
   if (k >= (1ULL << o.size())) return {};</pre>
   for (int i = int(o.size()) - 1; i >= 0; i--)
    if ((k >> i & 1) ^ (v >> o[i].second & 1))
     v ^= o[i].first;
   return v;
  Basis filter(S l) {
  Basis res = *this;
   for (int i = 0; i < BITS; i++)</pre>
   if (res.b[i].second < l) res.b[i] = {0, MIN};</pre>
   return res;
};
```

2.6 Binary Search On Segment Tree

```
find_first = x \rightarrow minimal x s.t. check([a, x))
// find_last = x \rightarrow maximal x s.t. check([x, b))
template <typename C>
int find_first(int l, const C &check) {
 if (l >= n) return n + 1;
 l += sz;
 for (int i = height; i > 0; i--)
  propagate(l >> i);
 Monoid sum = identity;
 do {
  while ((l & 1) == 0) l >>= 1;
  if (check(f(sum, data[l]))) {
   while (l < sz) {</pre>
    propagate(l);
    .
l <<= 1;
    auto nxt = f(sum, data[l]);
    if (not check(nxt)) {
     sum = nxt:
     l++;
    }
   }
   return l + 1 - sz;
  sum = f(sum, data[l++]);
 } while ((l & -l) != l);
 return n + 1:
template <typename C>
int find_last(int r, const C &check) {
 if (r <= 0) return -1;
 r += sz;
 for (int i = height; i > 0; i--)
  propagate((r - 1) >> i);
 Monoid sum = identity;
 do {
  while (r > 1 and (r & 1)) r >>= 1;
  if (check(f(data[r], sum))) {
   while (r < sz) {</pre>
    propagate(r);
    r = (r << 1) + 1;
    auto nxt = f(data[r], sum);
    if (not check(nxt)) {
     sum = nxt;
     r--;
    }
   return r - sz;
 sum = f(data[r], sum);
} while ((r & -r) != r);
 return -1;
```

3 Graph3.1 2-SAT (SCC)

```
class TwoSat{
```

```
private:
int n;
vector<vector<int>> rG,G,sccs;
vector<int> ord,idx;
vector<bool> vis,result;
void dfs(int u){
 vis[u]=true;
 for(int v:G[u])
   if(!vis[v]) dfs(v);
 ord.push_back(u);
void rdfs(int u){
 vis[u]=false;idx[u]=sccs.size()-1;
  sccs.back().push_back(u);
  for(int v:rG[u])
   if(vis[v])rdfs(v);
public:
void init(int n_){
 G.clear();G.resize(n=n_);
 rG.clear();rG.resize(n);
  sccs.clear();ord.clear();
  idx.resize(n);result.resize(n);
void add_edge(int u,int v){
 G[u].push_back(v);rG[v].push_back(u);
void orr(int x,int y){
 if ((x^y)==1)return;
  add_edge(x^1,y); add_edge(y^1,x);
bool solve(){
 vis.clear();vis.resize(n);
 for(int i=0;i<n;++i)</pre>
  if(not vis[i])dfs(i);
  reverse(ord.begin(),ord.end());
 for (int u:ord){
  if(!vis[u])continue;
   sccs.push_back(vector<int>());
  rdfs(u);
 for(int i=0;i<n;i+=2)</pre>
  if(idx[i]==idx[i+1])
   return false;
  vector<bool> c(sccs.size());
  for(size_t i=0;i<sccs.size();++i){</pre>
   for(auto sij : sccs[i]){
   result[sij]=c[i];
    c[idx[sij^1]]=!c[i];
  }
 }
  return true;
bool get(int x){return result[x];}
 int get_id(int x){return idx[x];}
int count(){return sccs.size();}
} sat2;
3.2 BCC
class BCC {
private:
```

```
int n, ecnt;
vector<vector<pair<int, int>>> g;
vector<int> dfn, low;
vector<bool> ap, bridge;
void dfs(int u, int f) {
 dfn[u] = low[u] = dfn[f] + 1;
 int ch = 0;
 for (auto [v, t] : g[u]) if (v != f) {
  if (dfn[v]) {
   low[u] = min(low[u], dfn[v]);
  } else {
   ++ch, dfs(v, u);
   low[u] = min(low[u], low[v]);
   if (low[v] > dfn[u])
    bridge[t] = true;
   if (low[v] >= dfn[u])
    ap[u] = true;
 }
 ap[u] &= (ch != 1 or u != f);
```

```
public:
    void init(int n_) {
        g.assign(n = n_, vector<pair<int, int>>());
        low.assign(n, ecnt = 0);
        dfn.assign(n, 0);
        ap.assign(n, false);
    }
    void add_edge(int u, int v) {
        g[u].emplace_back(v, ecnt);
        g[v].emplace_back(u, ecnt++);
    }
    void solve() {
        bridge.assign(ecnt, false);
        for (int i = 0; i < n; ++i)
              if (not dfn[i]) dfs(i, i);
    }
    bool is_ap(int x) { return ap[x]; }
    bool is_bridge(int x) { return bridge[x]; }
};</pre>
```

3.3 Round Square Tree

```
int N, M, cnt;
std::vector<int> G[maxn], T[maxn * 2];
int dfn[maxn], low[maxn], dfc;
int stk[maxn], tp;
void Tarjan(int u) {
low[u] = dfn[u] = ++dfc;
 stk[++tp] = u;
 for (int v : G[u]) {
  if (!dfn[v]) {
   Tarjan(v);
   low[u] = std::min(low[u], low[v]);
   if (low[v] == dfn[u]) {
    for (int x = 0; x != v; --tp) {
     x = stk[tp];
     T[cnt].push_back(x);
     T[x].push_back(cnt);
    T[cnt].push_back(u);
    T[u].push_back(cnt);
  } else
   low[u] = std::min(low[u], dfn[v]);
}
int main() { // ...
 cnt = N;
 for (int u = 1; u <= N; ++u)</pre>
  if (!dfn[u]) Tarjan(u), --tp;
```

3.4 Centroid Decomposition

```
struct Centroid {
 using G = vector<vector<pair<int, int>>>;
 vector<vector<int64_t>> Dist;
 vector<int> Pa, Dep;
 vector<int64_t> Sub, Sub2;
 vector<int> Cnt, Cnt2;
 vector<int> vis, sz, mx, tmp;
 void DfsSz(const G &g, int x) {
  vis[x] = true, sz[x] = 1, mx[x] = 0;
  for (auto [u, w] : g[x]) if (not vis[u]) {
   DfsSz(g, u); sz[x] += sz[u];
   mx[x] = max(mx[x], sz[u]);
  tmp.push_back(x);
 void DfsDist(const G &g, int x, int64_t D = 0) {
  Dist[x].push_back(D); vis[x] = true;
  for (auto [u, w] : g[x])
   if (not vis[u]) DfsDist(g, u, D + w);
 void DfsCen(const G &g, int x, int D = 0, int p = -1)
  tmp.clear(); DfsSz(g, x);
  int M = tmp.size(), C = -1;
  for (int u : tmp) {
   if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
```

```
vis[u] = false;
 DfsDist(g, C);
  for (int u : tmp) vis[u] = false;
  Pa[C] = p, vis[C] = true, Dep[C] = D;
  for (auto [u, w] : g[C])
  if (not vis[u]) DfsCen(g, u, D + 1, C);
Centroid(int N, G g)
   : Sub(N), Sub2(N), Cnt(N), Cnt2(N), Dist(N), Pa(N),
    Dep(N), vis(N), sz(N), mx(N) { DfsCen(g, 0); }
void Mark(int v) {
  int x = v, z = -1;
  for (int i = Dep[v]; i >= 0; --i) {
  Sub[x] += Dist[v][i], Cnt[x]++;
   if (z != -1)
   Sub2[z] += Dist[v][i], Cnt2[z]++;
   x = Pa[z = x];
 }
int64_t Query(int v) {
 int64_t res = 0;
 int x = v, z = -1;
for (int i = Dep[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Cnt[x] * Dist[v][i];
  if (z != -1)
   res -= Sub2[z] + 1LL * Cnt2[z] * Dist[v][i];
   x = Pa[z = x];
 }
  return res;
};
```

3.5 Directed Minimum Spanning Tree

```
struct Edge { int u, v, w; };
struct DirectedMST { // find maximum
 int solve(vector<Edge> E, int root, int n) {
  int ans = 0;
  while (true) {
   // find best in edge
   vector<int> in(n, -inf), prv(n, -1);
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
     in[e.v] = e.w;
     prv[e.v] = e.u;
   in[root] = 0; prv[root] = -1;
   for (int i = 0; i < n; i++)</pre>
    if (in[i] == -inf) return -inf;
   // find cycle
   int tot = 0;
   vector<int> id(n, -1), vis(n, -1);
   for (int i = 0; i < n; i++) {</pre>
    ans += in[i];
    for (int x = i; x != -1 && id[x] == -1; x = prv[x])
     if (vis[x] == i) {
      for (int y = prv[x]; y != x; y = prv[y])
       id[y] = tot;
      id[x] = tot++;
      break;
     vis[x] = i;
    }
   if (!tot) return ans;
   for (int i = 0; i < n; i++)</pre>
    if (id[i] == -1) id[i] = tot++;
   for (auto &e : E) {
    if (id[e.u] != id[e.v]) e.w -= in[e.v];
    e.u = id[e.u], e.v = id[e.v];
   n = tot; root = id[root];
  }
} DMST;
```

3.6 Dominator Tree

```
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
```

```
void init(int n) {
 // vertices are numbered from 0 to n - 1
 fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
 fill(fa, fa + n, -1); fill(val, val + n, -1);
 fill(sdom, sdom + n, -1);fill(rp, rp + n, -1);
 fill(dom, dom + n, -1); tk = 0;
for (int i = 0; i < n; ++i) {
  g[i].clear(); r[i].clear(); rdom[i].clear();
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
 rev[dfn[x] = tk] = x;
 fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
 for (int u : g[x]) {
 if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
  r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
 if (fa[x] == x) return c ? -1 : x;
 int p = find(fa[x], 1);
 if (p == -1) return c ? fa[x] : val[x];
 if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
 fa[x] = p;
 return c ? p : val[x];
vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
// p[i] = -2 if i is unreachable from s
 dfs(s);
 for (int i = tk - 1; i >= 0; --i) {
  for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
  if (i) rdom[sdom[i]].push_back(i);
  for (int &u : rdom[i]) {
   int p = find(u);
   if (sdom[p] == i) dom[u] = i;
   else dom[u] = p;
  if (i) merge(i, rp[i]);
 vector<int> p(n, -2); p[s] = -1;
 for (int i = 1; i < tk; ++i)</pre>
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
 for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
 return p;
}}
      Edge Coloring
```

```
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
 for (int i = 0; i <= N; i++)</pre>
  for (int j = 0; j <= N; j++)</pre>
    C[i][j] = G[i][j] = 0;
void solve(vector<pair<int, int>> &E, int N) {
 int X[kN] = {}, a;
 auto update = [&](int u) {
  for (X[u] = 1; C[u][X[u]]; X[u]++);
 auto color = [&](int u, int v, int c) {
  int p = G[u][v];
  G[u][v] = G[v][u] = c;
  C[u][c] = v, C[v][c] = u;
  C[u][p] = C[v][p] = 0;
  if (p) X[u] = X[v] = p;
  else update(u), update(v);
  return p;
 auto flip = [&](int u, int c1, int c2) {
  int p = C[u][c1];
  swap(C[u][c1], C[u][c2]);
  if (p) G[u][p] = G[p][u] = c2;
  if (!C[u][c1]) X[u] = c1;
  if (!C[u][c2]) X[u] = c2;
  return p;
 for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {</pre>
 auto [u, v] = E[t];
```

```
int v0 = v, c = X[u], c0 = c, d;
 vector<pair<int, int>> L; int vst[kN] = {};
 while (!G[u][v0]) {
  L.emplace_back(v, d = X[v]);
  if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
    c = color(u, L[a].first, c);
  else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
    color(u, L[a].first, L[a].second);
  else if (vst[d]) break;
  else vst[d] = 1, v = C[u][d];
 if (!G[u][v0]) {
  for (; v; v = flip(v, c, d), swap(c, d));
  if (C[u][c0]) { a = int(L.size()) - 1;
   while (--a >= 0 && L[a].second != c);
   for(;a>=0;a--)color(u,L[a].first,L[a].second);
  } else t--;
}
```

3.8 Lowbit Decomposition

```
class LBD {
int timer, chains;
vector<vector<int>> G;
vector<int> tl, tr, chain, head, dep, pa;
// chains : number of chain
// tl, tr[u] : subtree interval in the seq. of u
// head[i] : head of the chain i
// chian[u] : chain id of the chain u is on
void predfs(int u, int f) {
 dep[u] = dep[pa[u] = f] + 1;
  for (int v : G[u]) if (v != f) {
  predfs(v, u);
   if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
    chain[u] = chain[v];
 if (chain[u] == 0) chain[u] = ++chains;
void dfschain(int u, int f) {
 tl[u] = timer++;
  if (head[chain[u]] == -1)
  head[chain[u]] = u;
  for (int v : G[u])
  if (v != f and chain[v] == chain[u])
    dfschain(v, u);
  for (int v : G[u])
   if (v != f and chain[v] != chain[u])
    dfschain(v, u);
  tr[u] = timer;
public:
LBD(int n) : timer(0), chains(0), G(n), tl(n), tr(n),
chain(n), head(n, -1), dep(n), pa(n) {}
void add_edge(int u, int v) {
 G[u].push_back(v); G[v].push_back(u);
void decompose() { predfs(0, 0); dfschain(0, 0); }
PII get_subtree(int u) { return {tl[u], tr[u]}; }
vector<PII> get_path(int u, int v) {
 vector<PII> res;
while (chain[u] != chain[v]) {
   if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
    swap(u, v);
   int s = head[chain[u]];
  res.emplace_back(tl[s], tl[u] + 1);
  u = pa[s];
 if (dep[u] < dep[v]) swap(u, v);</pre>
 res.emplace_back(tl[v], tl[u] + 1);
  return res;
```

3.9 Manhattan Minimum Spanning Tree

```
typedef Point<int> P;
vector<array<int, 3>> manhattanMST(vector<P> ps) {
  vi id(sz(ps));
  iota(all(id), 0);
  vector<array<int, 3>> edges;
  rep(k, 0, 4) {
    sort(all(id), [&](int i, int j) {
```

```
return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;
});
map<int, int> sweep;
for (int i : id) {
   for (auto it = sweep.lower_bound(-ps[i].y);
        it != sweep.end(); sweep.erase(it++)) {
        int j = it->second;
        P d = ps[i] - ps[j];
        if (d.y > d.x) break;
        edges.push_back({d.y + d.x, i, j});
     }
     sweep[-ps[i].y] = i;
}
for (P &p : ps)
     if (k & 1) p.x = -p.x;
     else swap(p.x, p.y);
}
return edges; // [{w, i, j}, ...]
}
```

```
3.10 MaxClique
// contain a self loop u to u, than u won't in clique
template < size_t MAXN >
class MaxClique{
private:
 using bits = bitset< MAXN >;
 bits popped, G[ MAXN ], ans;
 size_t deg[ MAXN ], deo[ MAXN ], n;
 void sort_by_degree() {
  popped.reset();
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
    deg[ i ] = G[ i ].count();
  for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;</pre>
    for ( size_t j = 0 ; j < n ; ++ j )</pre>
      if ( not popped[ j ] and deg[ j ] < mi )
    mi = deg[ id = j ];</pre>
    popped[ deo[ i ] = id ] = 1;
    for( size_t u = G[ i ]._Find_first() ;
     u < n ; u = G[ i ]._Find_next( u ) )
       -- deg[ u ];
 void BK( bits R, bits P, bits X ) {
  if (R.count()+P.count() <= ans.count()) return;</pre>
  if ( not P.count() and not X.count() ) {
   if ( R.count() > ans.count() ) ans = R;
   return;
  /* greedily chosse max degree as pivot
  bits cur = P \mid X; size_t pivot = 0, sz = 0;
  for ( size_t u = cur._Find_first() ;
   u < n ; u = cur.\_Find\_next(u)
    if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
  cur = P & ( ~G[ pivot ] );
  */ // or simply choose first
  bits cur = P & (~G[ ( P | X )._Find_first() ]);
  for ( size_t u = cur._Find_first()
   u < n ; u = cur._Find_next( u ) ) {</pre>
   if ( R[ u ] ) continue;
   R[u] = 1;
   BK(R, P & G[u], X & G[u]);
   R[u] = P[u] = 0, X[u] = 1;
  }
public:
 void init( size_t n_ ) {
  n = n_{\cdot};
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
   G[ i ].reset();
  ans.reset();
 void add_edges( int u, bits S ) { G[ u ] = S; }
void add_edge( int u, int v ) {
  G[u][v] = G[v][u] = 1;
 int solve() {
  sort_by_degree(); // or simply iota( deo... )
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
   deg[ i ] = G[ i ].count();
  bits pob, nob = 0; pob.set();
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
```

```
for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t v = deo[ i ];
    bits tmp; tmp[ v ] = 1;
    BK( tmp, pob & G[ v ], nob & G[ v ] );
    pob[ v ] = 0, nob[ v ] = 1;
    }
    return static_cast< int >( ans.count() );
}
};
```

3.11 Minimum Mean Cycle

```
/* minimum mean cycle O(VE) */
struct MMC{
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
#define V 1021
#define inf 1e9
 struct Edge { int v,u; double c; };
 int n, m, prv[V][V], prve[V][V], vst[V];
 Edge e[E];
 vector<int> edgeID, cycle, rho;
 double d[V][V];
 void init( int _n ) { n = _n; m = 0; }
 // WARNING: TYPE matters
 void add_edge( int vi , int ui , double ci )
 { e[ m ++ ] = { vi , ui , ci }; }
 void bellman_ford() {
  for(int i=0; i<n; i++) d[0][i]=0;</pre>
  for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);</pre>
   for(int j=0; j<m; j++) {</pre>
    int v = e[j].v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
     d[i+1][u] = d[i][v]+e[j].c;
     prv[i+1][u] = v;
     prve[i+1][u] = j;
   }
  }
 double solve(){
  // returns inf if no cycle, mmc otherwise
  double mmc=inf;
  int st = -1;
  bellman_ford();
  for(int i=0; i<n; i++) {</pre>
   double avg=-inf;
   for(int k=0; k<n; k++) {</pre>
    if(d[n][i]<inf-eps)</pre>
     avg=max(avg,(d[n][i]-d[k][i])/(n-k));
    else avg=max(avg,inf);
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
  for (int i=n; !vst[st]; st=prv[i--][st]) {
   vst[st]++;
   edgeID.PB(prve[i][st]);
   rho.PB(st);
  while (vst[st] != 2) {
   int v = rho.back(); rho.pop_back();
   cycle.PB(v);
   vst[v]++;
  reverse(ALL(edgeID));
  edgeID.resize(SZ(cycle));
  return mmc;
} mmc;
```

3.12 Mo's Algorithm on Tree

```
dfs u:
  push u
  iterate subtree
  push u
Let P = LCA(u, v) with St(u)<=St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]</pre>
```

3.13 Virtual Tree

```
vector<pair<int, int>> build(vector<int> vs, int r) {
 vector<pair<int, int>> res;
 sort(vs.begin(), vs.end(), [](int i, int j) {
  return dfn[i] < dfn[j]; });</pre>
 vector<int> s = {r};
 for (int v : vs) if (v != r) {
  if (int o = lca(v, s.back()); o != s.back()) {
   while (s.size() >= 2) {
    if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
    res.emplace_back(s[s.size() - 2], s.back());
    s.pop_back();
   if (s.back() != o) {
    res.emplace_back(o, s.back());
    s.back() = o;
  s.push_back(v);
 for (size_t i = 1; i < s.size(); ++i)</pre>
  res.emplace_back(s[i - 1], s[i]);
 return res; // (x, y): x->y
```

4 Matching & Flow

4.1 Bipartite Matching

```
struct BipartiteMatching {
 vector<int> X[N];
 int fX[N], fY[N], n;
 bitset<N> vis;
 bool dfs(int x) {
  for (auto i : X[x]) if (not vis[i]) {
   vis[i] = true;
   if (fY[i] == -1 || dfs(fY[i])) {
    fY[fX[x] = i] = x;
    return true:
   }
  return false;
 void init(int n_, int m) {
  fill_n(X, n = n_, vector<int>());
  memset(fX, -1, sizeof(int) * n);
memset(fY, -1, sizeof(int) * m);
 void add_edge(int x, int y) { X[x].push_back(y); }
 int solve() { // return how many pair matched
  int cnt = 0;
  for (int i = 0; i < n; i++) {
   vis.reset();
   cnt += dfs(i);
  return cnt;
};
```

4.2 Dijkstra Cost Flow

```
// kN = #(vertices)
// MCMF.{Init, AddEdge, MincostMaxflow}
// MincostMaxflow(source, sink, flow_limit, &cost)
// => flow
using Pii = pair<int, int>;
constexpr int kInf = 0x3f3f3f3f, kN = 500;
struct Edge {
 int to, rev, cost, flow;
struct MCMF { // 0-based
int n{}, m{}, s{}, t{};
 vector<Edge> graph[kN];
 // Larger range for relabeling
 int64_t dis[kN] = {}, h[kN] = {};
 int p[kN] = {};
 void Init(int nn) {
 n = nn;
  for (int i = 0; i < n; i++) graph[i].clear();</pre>
 void AddEdge(int u, int v, int f, int c) {
  graph[u].push_back({v,
   static_cast<int>(graph[v].size()), c, f});
```

```
graph[v].push_back(
   {u, static_cast<int>(graph[u].size()) - 1,
    -c, 0});
bool Dijkstra(int &max_flow, int64_t &cost) {
  priority_queue<Pii, vector<Pii>, greater<>> pq;
  fill_n(dis, n, kInf);
  dis[s] = 0;
 pq.emplace(0, s);
 while (!pq.empty()) {
  auto u = pq.top();
   pq.pop();
   int v = u.second;
   if (dis[v] < u.first) continue;</pre>
   for (auto &e : graph[v]) {
    auto new_dis =
     dis[v] + e.cost + h[v] - h[e.to];
    if (e.flow > 0 && dis[e.to] > new_dis) {
     dis[e.to] = new_dis;
     p[e.to] = e.rev;
     pq.emplace(dis[e.to], e.to);
  if (dis[t] == kInf) return false;
  for (int i = 0; i < n; i++) h[i] += dis[i];</pre>
  int d = max_flow;
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   d = min(d, graph[e.to][e.rev].flow);
  }
 max_flow -= d;
 cost += int64_t(d) * h[t];
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   e.flow += d;
  graph[e.to][e.rev].flow -= d;
  return true;
 int MincostMaxflow(
 int ss, int tt, int max_flow, int64_t &cost) {
 this->s = ss, this->t = tt;
  cost = 0;
 fill_n(h, n, 0);
 auto orig_max_flow = max_flow;
 while (Dijkstra(max_flow, cost) && max_flow) {}
  return orig_max_flow - max_flow;
};
```

4.3 Dinic

```
template <typename Cap = int64_t>
class Dinic{
private:
 struct E{
    int to, rev;
    Cap cap;
 int n, st, ed;
 vector<vector<E>> G;
 vector<int> lv, idx;
 bool BFS(){
   lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
   while (not bfs.empty()){
      int u = bfs.front(); bfs.pop();
      for (auto e: G[u]) {
        if (e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
        bfs.push(e.to); lv[e.to] = lv[u] + 1;
      }
    return lv[ed] != -1;
 Cap DFS(int u, Cap f){
   if (u == ed) return f;
    Cap ret = 0;
    for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
      auto &e = G[u][i];
```

```
if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
        Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
        G[e.to][e.rev].cap += nf;
        if (f == 0) return ret:
      if (ret == 0) lv[u] = -1;
     return ret;
public:
   void init(int n_) { G.assign(n = n_, vector<E>()); }
   void add_edge(int u, int v, Cap c){
     G[u].push_back({v, int(G[v].size()), c});
G[v].push_back({u, int(G[u].size())-1, 0});
   Cap max_flow(int st_, int ed_){
  st = st_, ed = ed_; Cap ret = 0;
     while (BFS()) {
        idx.assign(n, 0);
        Cap f = DFS(st, numeric_limits<Cap>::max());
        ret += f;
        if (f == 0) break;
      return ret;
   }
};
```

4.4 Flow Models

- · Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.

 - 2. For each edge (x,y,l,u), connect $x\to y$ with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, con- $\text{nect } v \to T \text{ with capacity } -in(v).$
 - To maximize, connect t
 ightarrow s with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Oth-
 - erwise, the maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to Tbe f'. If $f + f' \neq \sum_{v \in V, in(v) > 0}^{\bullet} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- · Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise. 2. DFS from unmatched vertices in X.

 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x, y, c), connect $x \to y$ with (cost, cap) = (c, 1) if c>0, otherwise connect $y\to x$ with (cost, cap)=(-c,1)3. For each edge with c<0, sum these cost as K, then increase
 - d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v) > 0, connect $S \rightarrow v$ with (cost, cap) =
 - 5. For each vertex v with d(v) < 0, connect v \rightarrow T with
 - (cost, cap) = (0, -d(v)) 6. Flow from S to T , the answer is the cost of the flow C+K
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\it T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \to v$, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
 - 5. For $v \in G$, connect it with sink $v \to t$ with capacity K + 2T t $(\sum_{e \in E(v)} w(e)) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u, v).
 - 2. Connect v o v' with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G^{\prime} .
- · Project selection problem
 - 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.

· 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with capacity c_y .
- 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4.5 General Graph Matching

```
namespace matching {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
 g[u].push_back(v);
 g[v].push_back(u);
int Find(int u) {
return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
 static int tk = 0; tk++;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
 }
void Blossom(int x, int y, int l) {
 while (Find(x) != l) {
  pre[x] = y, y = match[x];
  if (s[y] = 1) q.push(y), s[y] = 0;
  if (fa[x] == x) fa[x] = l;
  if (fa[y] == y) fa[y] = l;
  x = pre[y];
bool Bfs(int r, int n) {
 for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;
 while (!q.empty()) q.pop();
 q.push(r);
 s[r] = 0;
 while (!q.empty()) {
  int x = q.front(); q.pop();
  for (int u : g[x]) {
   if (s[u] == -1) {
    pre[u] = x, s[u] = 1;
    if (match[u] == n) {
     for (int a = u, b = x, last; b != n; a = last, b =
     pre[a])
      last = match[b], match[b] = a, match[a] = b;
     return true;
    q.push(match[u]);
    s[match[u]] = 0;
   } else if (!s[u] && Find(u) != Find(x)) {
    int l = LCA(u, x, n);
    Blossom(x, u, l);
    Blossom(u, x, l);
  }
 return false;
int Solve(int n) {
 int res = 0;
 for (int x = 0; x < n; ++x) {
  if (match[x] == n) res += Bfs(x, n);
 return res;
```

4.6 Global Min-Cut

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
pair<int, int> phase(int n) {
 memset(v, false, sizeof(v));
 memset(g, 0, sizeof(g));
 int s = -1, t = -1;
 while (true) {
  int c = -1;
  for (int i = 0; i < n; ++i) {</pre>
   if (del[i] || v[i]) continue;
   if (c == -1 || g[i] > g[c]) c = i;
  if (c == -1) break;
  v[s = t, t = c] = true;
  for (int i = 0; i < n; ++i) {</pre>
   if (del[i] || v[i]) continue;
   g[i] += w[c][i];
 }
 return make_pair(s, t);
int mincut(int n) {
 int cut = 1e9;
 memset(del, false, sizeof(del));
 for (int i = 0; i < n - 1; ++i) {</pre>
 int s, t; tie(s, t) = phase(n);
del[t] = true; cut = min(cut, g[t]);
  for (int j = 0; j < n; ++j) {</pre>
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
  }
 }
 return cut;
```

4.7 GomoryHu Tree

```
int g[maxn];
vector<edge> GomoryHu(int n){
 vector<edge> rt;
 for(int i=1;i<=n;++i)g[i]=1;</pre>
 for(int i=2;i<=n;++i){</pre>
  int t=g[i];
  flow.reset(); // clear flows on all edge
  rt.push_back({i,t,flow(i,t)});
  flow.walk(i); // bfs points that connected to i (use
    edges not fully flow)
  for(int j=i+1;j<=n;++j){</pre>
   if(g[j]==t && flow.connect(j))g[j]=i; // check if i
    can reach i
  }
 }
 return rt;
```

4.8 Kuhn Munkres

lld d;

```
class KM {
private:
 static constexpr lld INF = 1LL << 60;</pre>
 vector<lld> hl,hr,slk;
 vector<int> fl,fr,pre,qu;
 vector<vector<lld>> w;
 vector<bool> vl,vr;
 int n, ql, qr;
 bool check(int x) {
  if (vl[x] = true, fl[x] != -1)
   return vr[qu[qr++] = fl[x]] = true;
  while (x != -1) swap(x, fr[fl[x] = pre[x]]);
  return false;
 void bfs(int s) {
  fill(slk.begin(), slk.end(), INF);
fill(vl.begin(), vl.end(), false);
fill(vr.begin(), vr.end(), false);
  ql = qr = 0;
  vr[qu[qr++] = s] = true;
  while (true) {
```

```
while (ql < qr) {</pre>
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
     if(!vl[x]&&slk[x]>=(d=hl[x]+hr[y]-w[x][y])){
      if (pre[x] = y, d) slk[x] = d;
      else if (!check(x)) return;
     }
   }
  d = INF;
  for (int x = 0; x < n; ++x)
   if (!vl[x] && d > slk[x]) d = slk[x];
   for (int x = 0; x < n; ++x) {
   if (vl[x]) hl[x] += d;
    else slk[x] -= d;
    if (vr[x]) hr[x] -= d;
   for (int x = 0; x < n; ++x)
    if (!vl[x] && !slk[x] && !check(x)) return;
 }
public:
void init( int n_ ) {
 qu.resize(n = n_);
 fl.assign(n, -1); fr.assign(n, -1);
 hr.assign(n, 0); hl.resize(n);
 w.assign(n, vector<lld>(n));
 slk.resize(n); pre.resize(n);
 vl.resize(n); vr.resize(n);
}
void set_edge( int u, int v, lld x ) {w[u][v] = x;}
lld solve() {
 for (int i = 0; i < n; ++i)</pre>
  hl[i] = *max_element(w[i].begin(), w[i].end());
 for (int i = 0; i < n; ++i) bfs(i);</pre>
 lld res = 0;
 for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
 return res:
}
} km:
```

4.9 Minimum Cost Circulation

```
struct Edge { int to, cap, rev, cost; };
vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
bool mark[kN];
int NegativeCycle(int n) {
memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
 int upd = -1;
for (int i = 0; i <= n; ++i) {
 for (int j = 0; j < n; ++j) {</pre>
   int idx = 0;
  for (auto &e : g[j]) {
    if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
     dist[e.to] = dist[j] + e.cost;
     pv[e.to] = j, ed[e.to] = idx;
     if (i == n) {
     upd = j;
     while(!mark[upd])mark[upd]=1,upd=pv[upd];
     return upd;
     }
    idx++;
  }
 }
return -1;
int Solve(int n) {
int rt = -1, ans = 0;
while ((rt = NegativeCycle(n)) >= 0) {
 memset(mark, false, sizeof(mark));
 vector<pair<int, int>> cyc;
 while (!mark[rt]) {
  cyc.emplace_back(pv[rt], ed[rt]);
  mark[rt] = true;
  rt = pv[rt];
 reverse(cyc.begin(), cyc.end());
 int cap = kInf;
  for (auto &i : cyc) {
  auto &e = g[i.first][i.second];
```

```
cap = min(cap, e.cap);
}
for (auto &i : cyc) {
  auto &e = g[i.first][i.second];
  e.cap -= cap;
  g[e.to][e.rev].cap += cap;
  ans += e.cost * cap;
}
}
return ans;
}
```

```
Minimum Cost Maximum Flow
4.10
template <typename Cap, typename Wei> class MCMF {
 static constexpr auto INF_CAP = numeric_limits<Cap>::
    max();
 static constexpr auto INF_WEI = numeric_limits<Wei>::
    max();
private:
 struct E {
  int to, rev;
  Cap cap; Wei wei;
  E() {}
  E(int a, int b, Cap c, Wei d) : to(a), rev(b), cap(c)
    , wei(d) {}
 };
 int S, T;
 vector<vector<E>> G;
 vector<pair<int, int>> f;
 vector<int> inq;
 vector<Wei> d; vector<Cap> up;
 optional<pair<Cap, Wei>> SPFA() {
  queue<int> q:
  for (q.push(S), d[S] = 0, up[S] = INF_CAP; not q.
    empty(); q.pop()) {
   int u = q.front(); inq[u] = false;
   if (up[u] == 0) continue;
   for (int i = 0; i < int(G[u].size()); ++i) {</pre>
    auto e = G[u][i]; int v = e.to;
    if (e.cap <= 0 or d[v] <= d[u] + e.wei)</pre>
    continue:
    d[v] = d[u] + e.wei; f[v] = {u, i};
    up[v] = min(up[u], e.cap);
    if (not inq[v]) q.push(v);
    inq[v] = true;
  if (d[T] == INF_WEI) return nullopt;
  for (int i = T; i != S; i = f[i].first) {
   auto &eg = G[f[i].first][f[i].second];
   eg.cap -= up[T];
   G[eg.to][eg.rev].cap += up[T];
  return pair{up[T], d[T]};
public:
 void init(int n) {
  G.assign(n, vector<E>());
  f.resize(n), up.resize(n);
  inq.assign(n, false), d.assign(n, INF_WEI);
 void add_edge(int s, int t, Cap c, Wei w) {
 G[s].emplace_back(t, int(G[t].size()), c, w);
  G[t].emplace_back(s, int(G[s].size()) - 1, 0, -w);
 pair<Cap, Wei> solve(int a, int b) {
  S = a, T = b;
  Cap c = 0; Wei w = 0;
  while (auto r = SPFA()) {
   c += r->first, w += r->first * r->second;
   fill(inq.begin(), inq.end(), false);
   fill(d.begin(), d.end(), INF_WEI);
  return {c, w};
 }
};
```

4.11 Maximum Weight Graph Matching

```
struct WeightGraph {
  static const int inf = INT_MAX;
```

```
static const int maxn = 514;
struct edge {
int u, v, w;
edge(){}
edge(int u, int v, int w): u(u), v(v), w(w) {}
int n, n_x;
edge g[maxn * 2][maxn * 2];
int lab[maxn * 2];
int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa
int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
   maxn * 2];
vector<int> flo[maxn * 2];
queue<int> q;
int e_delta(const edge &e) { return lab[e.u] + lab[e.v
   ] - g[e.u][e.v].w * 2; }
void update_slack(int u, int x) { if (!slack[x] ||
   e_delta(g[u][x]) < e_delta(g[slack[x]][x])) slack[x</pre>
   ] = u; }
void set_slack(int x) {
 slack[x] = 0;
 for (int u = 1; u <= n; ++u)</pre>
  if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
   update_slack(u, x);
void q_push(int x) {
 if (x <= n) q.push(x);
 else for (size_t i = 0; i < flo[x].size(); i++)</pre>
   q_push(flo[x][i]);
void set_st(int x, int b) {
 st[x] = b;
 if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)</pre>
    set_st(flo[x][i], b);
int get_pr(int b, int xr) {
 int pr = find(flo[b].begin(), flo[b].end(), xr) - flo
   [b].begin();
if (pr % 2 == 1) {
  reverse(flo[b].begin() + 1, flo[b].end());
 return (int)flo[b].size() - pr;
 return pr;
}
void set_match(int u, int v) {
match[u] = g[u][v].v;
if (u <= n) return;</pre>
 edge e = g[u][v];
int xr = flo_from[u][e.u], pr = get_pr(u, xr);
for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo</pre>
   [u][i ^ 1]);
set_match(xr, v);
rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
void augment(int u, int v) {
for (; ; ) {
  int xnv = st[match[u]];
  set_match(u, v);
  if (!xnv) return;
  set_match(xnv, st[pa[xnv]]);
  u = st[pa[xnv]], v = xnv;
}
int get_lca(int u, int v) {
static int t = 0;
 for (++t; u || v; swap(u, v)) {
 if (u == 0) continue;
 if (vis[u] == t) return u;
  vis[u] = t;
 u = st[match[u]];
  if (u) u = st[pa[u]];
 return 0;
void add_blossom(int u, int lca, int v) {
int b = n + 1;
 while (b <= n_x && st[b]) ++b;</pre>
 if (b > n_x) ++n_x;
 lab[b] = 0, S[b] = 0;
match[b] = match[lca];
```

```
flo[b].clear();
 flo[b].push_back(lca);
 for (int x = u, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
 reverse(flo[b].begin() + 1, flo[b].end());
 for (int x = v, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
 set_st(b, b);
 for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].w
   = 0;
 for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  for (int x = 1; x <= n_x; ++x)
   if (g[b][x].w == 0 || e_delta(g[xs][x]) < e_delta(g</pre>
   [b][x])
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
  for (int x = 1; x <= n; ++x)
   if (flo_from[xs][x]) flo_from[b][x] = xs;
 set_slack(b);
void expand_blossom(int b) {
for (size_t i = 0; i < flo[b].size(); ++i)
    set_st(flo[b][i], flo[b][i]);</pre>
 int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
   xr);
 for (int i = 0; i < pr; i += 2) {
  int xs = flo[b][i], xns = flo[b][i + 1];
  pa[xs] = g[xns][xs].u;
  S[xs] = 1, S[xns] = 0;
slack[xs] = 0, set_slack(xns);
  q_push(xns);
 S[xr] = 1, pa[xr] = pa[b];
 for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  S[xs] = -1, set_slack(xs);
st[b] = 0;
bool on_found_edge(const edge &e) {
 int u = st[e.u], v = st[e.v];
 if (S[v] == -1) {
 pa[v] = e.u, S[v] = 1;
  int nu = st[match[v]];
  slack[v] = slack[nu] = 0;
  S[nu] = 0, q_push(nu);
 } else if (S[v] == 0) {
  int lca = get_lca(u, v);
  if (!lca) return augment(u,v), augment(v,u), true;
  else add_blossom(u, lca, v);
 return false;
bool matching() {
memset(S + 1, -1, sizeof(int) * n_x);
 memset(slack + 1, 0, sizeof(int) * n_x);
 q = queue<int>();
 for (int x = 1; x <= n_x; ++x)</pre>
  if (st[x] == x \&\& !match[x]) pa[x] = 0, S[x] = 0,
   q_push(x);
 if (q.empty()) return false;
 for (; ; ) {
  while (q.size()) {
   int u = q.front(); q.pop();
   if (S[st[u]] == 1) continue;
   for (int v = 1; v <= n; ++v)</pre>
    if (g[u][v].w > 0 && st[u] != st[v]) {
    if (e_delta(g[u][v]) == 0) {
      if (on_found_edge(g[u][v])) return true;
     } else update_slack(u, st[v]);
  int d = inf;
  for (int b = n + 1; b <= n_x; ++b)
   if (st[b] == b && S[b] == 1) d = min(d, lab[b] / 2)
  for (int x = 1; x <= n_x; ++x)</pre>
   if (st[x] == x && slack[x]) {
```

```
if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x])
     else if (S[x] == 0) d = min(d, e_delta(g[slack[x
    ]][x]) / 2);
   for (int u = 1; u <= n; ++u) {</pre>
    if (S[st[u]] == 0) {
     if (lab[u] <= d) return 0;</pre>
     lab[u] -= d;
    } else if (S[st[u]] == 1) lab[u] += d;
   for (int b = n + 1; b \le n_x; ++b)
    if (st[b] == b) {
     if (S[st[b]] == 0) lab[b] += d * 2;
     else if (S[st[b]] == 1) lab[b] -= d * 2;
   q = queue<int>();
   for (int x = 1; x <= n_x; ++x)
    if (st[x] == x && slack[x] && st[slack[x]] != x &&
    e_delta(g[slack[x]][x]) == 0)
     if (on_found_edge(g[slack[x]][x])) return true;
   for (int b = n + 1; b \le n_x; ++b)
    if (st[b] == b && S[b] == 1 && lab[b] == 0)
    expand_blossom(b);
  return false;
 }
 pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n_x = n;
  int n_matches = 0;
  long long tot_weight = 0;
  for (int u = 0; u <= n; ++u) st[u] = u, flo[u].clear</pre>
    ();
  int w_max = 0;
  for (int u = 1; u <= n; ++u)</pre>
   for (int v = 1; v <= n; ++v) {</pre>
    flo_from[u][v] = (u == v ? u : 0);
    w_max = max(w_max, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
  while (matching()) ++n_matches;
  for (int u = 1; u <= n; ++u)
   if (match[u] && match[u] < u)</pre>
    tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
 void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
    g[vi][ui].w = wi; }
 void init(int _n) {
  n = _n;
  for (int u = 1; u <= n; ++u)</pre>
   for (int v = 1; v <= n; ++v)</pre>
    g[u][v] = edge(u, v, 0);
 }
};
```

5 Math

5.1 Common Bounds

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \backslash \{0\}} (-1)^{k+1} p(n - k(3k-1)/2)$$

$$p(n) \approx 0.145/n \cdot \exp(2.56\sqrt{n})$$

$$\frac{n}{\max_{i \le n} (d(i))} \frac{100 \text{ le3 le6 le9 lel2 lel5 lel8}}{12 \text{ 32 240 l344 6720 26880 l03680}}$$

$$\frac{n}{\binom{2n}{n}} \frac{12 \text{ 3 4 5 6 7 8 9 l0}}{2 \text{ 6 20 70 252 924 3432 l2870 48620 l84756}}$$

5.2 Strling Number

5.2.1 First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$
$$x(x+1) \dots (x+n-1) = \sum_{k=0}^{n} S_1(n,k) x^k$$

$$g(x) = x(x+1)\dots(x+n-1) = \sum_{k=0}^{n} a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^{n} \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^{k} ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.2.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k \binom{k}{i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b==(x/y)*a;
}
```

5.4 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
 vector<T> d(output.size() + 1), me, he;
 for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
  for (size_t j = 0; j < me.size(); ++j)</pre>
   d[i] += output[i - j - 2] * me[j];
  if ((d[i] -= output[i - 1]) == 0) continue;
  if (me.empty()) {
   me.resize(f = i);
   continue:
  vector<T> o(i - f - 1);
  T k = -d[i] / d[f]; o.push_back(-k);
  for (T x : he) o.push_back(x * k);
  if (o.size() < me.size()) o.resize(me.size());</pre>
  for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
  if (i-f+he.size() >= me.size()) he = me, f = i;
  me = o;
 return me;
```

5.5 Charateristic Polynomial

int N = A.size();

```
vector<vector<int>>> Hessenberg(const vector<vector<int</pre>
    >> &A) {
 int N = A.size();
 vector<vector<int>> H = A;
 for (int i = 0; i < N - 2; ++i) {
  if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {</pre>
    if (H[j][i]) {
     for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j</pre>
    ][k]);
     for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k</pre>
    ][j]);
     break;
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {</pre>
   int coef = 1LL * val * H[j][i] % kP;
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
     * H[i + 1][k] * (kP - coef)) % kP;
   for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
 }
 return H;
vector<int> CharacteristicPoly(const vector<vector<int</pre>
    >> &A) {
```

```
auto H = Hessenberg(A);
for (int i = 0; i < N; ++i) {</pre>
 for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];</pre>
vector<vector<int>> P(N + 1, vector<int>(N + 1));
P[0][0] = 1;
for (int i = 1; i <= N; ++i) {</pre>
P[i][0] = 0;
 for (int j = 1; j <= i; ++j) P[i][j] = P[i - 1][j -</pre>
   17:
 int val = 1;
 for (int j = i - 1; j >= 0; --j) {
  int coef = 1LL * val * H[j][i - 1] % kP;
  for (int k = 0; k <= j; ++k) P[i][k] = (P[i][k] + 1</pre>
   LL * P[j][k] * coef) % kP;
  if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 }
if (N & 1) {
 for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
return P[N];
```

5.6 Chinese Remainder

```
x = a1 % m1
x = a2 % m2
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
// 0 <= x < lcm(m1, m2)</pre>
```

5.7 De-Bruijn

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
 if (n % p == 0)
   for (int i = 1; i <= p; ++i)</pre>
    res[sz++] = aux[i];
 aux[t] = aux[t - p];
 db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {</pre>
  aux[t] = i;
   db(t + 1, t, n, k);
 }
}
int de_bruijn(int k, int n) {
// return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
if (k == 1) {
 res[0] = 0;
 return 1;
for (int i = 0; i < k * n; i++) aux[i] = 0;</pre>
sz = 0;
db(1, 1, n, k);
return sz;
}
```

5.8 DiscreteLog

```
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
    // x^? \equiv y (mod M)
    Int t = 1, c = 0, g = 1;
    for (Int M_ = M; M_ > 0; M_ >>= 1)
        g = g * x % M;
    for (g = gcd(g, M); t % g != 0; ++c) {
        if (t == y) return c;
        t = t * x % M;
    }
    if (y % g != 0) return -1;
    t /= g, y /= g, M /= g;
    Int h = 0, gs = 1;
    for (; h * h < M; ++h) gs = gs * x % M;
    unordered_map<Int, Int> bs;
    for (Int s = 0; s < h; bs[y] = ++s)
        y = y * x % M;
    for (Int s = 0; s < M; s += h) {</pre>
```

```
t = t * gs % M;
if (bs.count(t)) return c + s + h - bs[t];
}
return -1;
}
```

5.9 Extended Euler

```
a^b \equiv \begin{cases} a^{(b \mod \varphi(m)) + \varphi(m)} & \text{if } (a,m) \neq 1 \wedge b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m
```

5.10 ExtendedFloorSum

```
g(a, b, c, n) = \sum_{i=0}^{n} i \lfloor \frac{ai + b}{c} \rfloor
                            \left( \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \right)
                             +g(a \bmod c, b \bmod c, c, n),
                                                                                                           a \geq c \vee b \geq c
                                                                                                           n < 0 \lor a = 0
                             \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
                           -h(c, c-b-1, a, m-1)),
                                                                                                           otherwise
h(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^2
                            \left( \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right)
                             +\lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1)
                            +h(a \bmod c, b \bmod c, c, n)
                            +2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)
                            +2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
                                                                                                          a > c \lor b > c
                            0,
                                                                                                          n < 0 \lor a = 0
                            nm(m+1) - 2g(c, c-b-1, a, m-1)
                            -2f(c, c-b-1, a, m-1) - f(a, b, c, n), otherwise
```

5.11 Fast Fourier Transform

```
const int mod = 10000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
  static_assert (M1 <= M2 && M2 <= M3);</pre>
  constexpr int64_t r12 = modpow(M1, M2-2, M2);
  constexpr int64_t r13 = modpow(M1, M3-2, M3);
  constexpr int64_t r23 = modpow(M2, M3-2, M3);
  constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
  B = (B - A + M2) * r12 % M2;

C = (C - A + M3) * r13 % M3;
  C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
 for (int i = 0; i <= maxn; i++)</pre>
  omega[i] = cplx(cos(2 * pi * j / maxn),
     sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
  int z = __builtin_ctz(n) - 1;
 for (int i = 0; i < n; ++i) {</pre>
  int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^=(i >> j & 1)<<(z - j);
  if (x > i) swap(v[x], v[i]);
 for (int s = 2; s <= n; s <<= 1) {
  int z = s >> 1;
  for (int i = 0; i < n; i += s) {</pre>
   for (int k = 0; k < z; ++k) {
    cplx x = v[i + z + k] * omega[maxn / s * k];
    v[i + z + k] = v[i + k] - x;
    v[i + k] = v[i + k] + x;
 }
void ifft(vector<cplx> &v, int n) {
 fft(v, n); reverse(v.begin() + 1, v.end());
 for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
```

if (a >= m) {

if (b >= m) {

swap(m, a);

ans += n * (n - 1) / 2 * (a / m); a %= m;

 $n = (llu)(y_max / m), b = (llu)(y_max % m);$

ans += n * (b / m); b %= m;

llu y_max = a * n + b;
if (y_max < m) break;</pre>

// y_max < m * (n + 1) // floor(y_max / m) <= n

```
VL convolution(const VI &a, const VI &b) {
                                                              return ans;
// Should be able to handle N <= 10^5, C <= 10^4
                                                             lld floor_sum(lld n, lld m, lld a, lld b) {
int sz = 1;
                                                              llu ans = 0;
while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
vector<cplx> v(sz);
                                                              if (a < 0) {
for (int i = 0; i < sz; ++i) {</pre>
                                                               llu a2 = (a \% m + m) \% m;
 double re = i < a.size() ? a[i] : 0;
                                                               ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
 double im = i < b.size() ? b[i] : 0;</pre>
  v[i] = cplx(re, im);
                                                              if (b < 0) {
fft(v, sz);
                                                               llu b2 = (b % m + m) % m;
for (int i = 0; i <= sz / 2; ++i) {
  int j = (sz - i) & (sz - 1);</pre>
                                                               ans -= 1ULL * n * ((b2 - b) / m);
                                                               b = b2:
 cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
    * cplx(0, -0.25);
                                                              return ans + floor_sum_unsigned(n, m, a, b);
 if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
    ].conj()) * cplx(0, -0.25);
                                                             5.13 FWT
 v[i] = x;
                                                             /* or convolution:
 ifft(v, sz);
                                                              * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
VL c(sz);
                                                              * and convolution:
for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
                                                                x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
                                                             void fwt(int x[], int N, bool inv = false) {
                                                               for (int d = 1; d < N; d <<= 1) {</pre>
VI convolution_mod(const VI &a, const VI &b, int p) {
                                                                 for (int s = 0, d2 = d * 2; s < N; s += d2)
int sz = 1;
                                                                   for (int i = s, j = s + d; i < s + d; i++, j++) {
while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
                                                                     int ta = x[i], tb = x[j];
vector<cplx> fa(sz), fb(sz);
                                                                     x[i] = modadd(ta, tb);
for (int i = 0; i < (int)a.size(); ++i)</pre>
                                                                     x[j] = modsub(ta, tb);
 fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
for (int i = 0; i < (int)b.size(); ++i)</pre>
  fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
                                                               if (inv) for (int i = 0, invn = modinv(N); i < N; i</pre>
 fft(fa, sz), fft(fb, sz);
double r = 0.25 / sz;
                                                                 x[i] = modmul(x[i], invn);
cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
  int j = (sz - i) & (sz - 1);
                                                             5.14 Miller Rabin
 cplx a1 = (fa[i] + fa[j].conj());
                                                             bool isprime(llu x) {
 cplx a2 = (fa[i] - fa[j].conj()) * r2;
                                                              static auto witn = [](llu a, llu n, int t) {
 cplx b1 = (fb[i] + fb[j].conj()) * r3;
                                                               if (!a) return false;
  cplx b2 = (fb[i] - fb[j].conj()) * r4;
                                                               while (t--) {
  if (i != j) {
  cplx c1 = (fa[j] + fa[i].conj());
                                                                llu a2 = mmul(a, a, n);
                                                                if (a2 == 1 && a != 1 && a != n - 1) return true;
  cplx c2 = (fa[j] - fa[i].conj()) * r2;
                                                                a = a2;
  cplx d1 = (fb[j] + fb[i].conj()) * r3;
   cplx d2 = (fb[j] - fb[i].conj()) * r4;
                                                               return a != 1;
   fa[i] = c1 * d1 + c2 * d2 * r5;
  fb[i] = c1 * d2 + c2 * d1;
                                                              if (x < 2) return false;</pre>
                                                              if (!(x & 1)) return x == 2;
  fa[j] = a1 * b1 + a2 * b2 * r5;
                                                              int t = __builtin_ctzll(x - 1);
 fb[j] = a1 * b2 + a2 * b1;
                                                              llu odd = (x - 1) >> t;
                                                              for (llu m:
fft(fa, sz), fft(fb, sz);
                                                               {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
vector<int> res(sz);
                                                               if (witn(mpow(m % x, odd, x), x, t))
 for (int i = 0; i < sz; ++i) {</pre>
                                                                return false;
 long long a = round(fa[i].re), b = round(fb[i].re),
                                                              return true;
       c = round(fa[i].im);
 res[i] = (a+((b % p) << 15)+((c % p) << 30)) % p;
}
                                                             5.15 NTT
return res;
                                                             template <int mod, int G, int maxn>
}}
                                                             struct NTT {
5.12 FloorSum
                                                              static_assert (maxn == (maxn & -maxn));
// @param n `n < 2^32`
                                                              int roots[maxn];
// @param m `1 <= m < 2^32`
                                                              NTT () {
// @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
                                                               int r = modpow(G, (mod - 1) / maxn);
llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
                                                               for (int i = maxn >> 1; i; i >>= 1) {
llu ans = 0:
                                                                roots[i] = 1;
                                                                for (int j = 1; j < i; j++)</pre>
while (true) {
```

roots[i + j] = modmul(roots[i + j - 1], r);

void operator()(int F[], int n, bool inv = false) {

// n must be 2^k , and $0 \le F[i] \le mod$

for (int i = 0, j = 0; i < n; i++) {</pre>

for (int k = n>>1; (j^=k) < k; k>>=1);

if (i < j) swap(F[i], F[j]);</pre>

for (int s = 1; s < n; s *= 2) {
 for (int i = 0; i < n; i += s * 2) {</pre>

for (int j = 0; j < s; j++) {

r = modmul(r, r);

}

```
int a = F[i+j];
     int b = modmul(F[i+j+s], roots[s+j]);
     F[i+j] = modadd(a, b); // a + b
     F[i+j+s] = modsub(a, b); // a - b
  }
 if (inv) {
  int invn = modinv(n);
  for (int i = 0; i < n; i++)</pre>
   F[i] = modmul(F[i], invn);
   reverse(F + 1, F + n);
 }
}
NTT<2013265921, 31, 1048576> ntt;
```

5.16 Partition Number

```
int b = sqrt(n);
ans[0] = tmp[0] = 1;
for (int i = 1; i <= b; i++) {</pre>
 for (int rep = 0; rep < 2; rep++)
for (int j = i; j <= n - i * i; j++)</pre>
 modadd(tmp[j], tmp[j-i]);
for (int j = i * i; j <= n; j++)</pre>
  modadd(ans[j], tmp[j - i * i]);
```

5.17 Pi Count (Linear Sieve)

```
static constexpr int N = 1000000 + 5;
lld pi[N];
vector<int> primes;
bool sieved[N];
lld cube_root(lld x){
lld s=cbrt(x-static_cast<long double>(0.1));
while(s*s*s <= x) ++s;
return s-1;
lld square_root(lld x){
lld s=sqrt(x-static_cast<long double>(0.1));
while(s*s <= x) ++s;
return s-1;
void init(){
primes.reserve(N);
primes.push_back(1);
for(int i=2;i<N;i++) {</pre>
 if(!sieved[i]) primes.push_back(i);
 pi[i] = !sieved[i] + pi[i-1];
 for(int p: primes) if(p > 1) {
   if(p * i >= N) break;
   sieved[p * i] = true;
   if(p % i == 0) break;
 }
}
lld phi(lld m, lld n) {
static constexpr int MM = 80000, NN = 500;
static lld val[MM][NN];
if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
 if(n == 0) return m;
if(primes[n] >= m) return 1;
lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
 if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
return ret;
lld pi_count(lld);
lld P2(lld m, lld n) {
lld sm = square_root(m), ret = 0;
for(lld i = n+1;primes[i]<=sm;i++)</pre>
 ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
return ret;
lld pi_count(lld m) {
if(m < N) return pi[m];</pre>
lld n = pi_count(cube_root(m));
return phi(m, n) + n - 1 - P2(m, n);
```

```
// does not work when n is prime
// return any non-trivial factor
llu pollard_rho(llu n) {
 static auto f = [](llu x, llu k, llu m) {
    return add(k, mul(x, x, m), m); };
 if (!(n & 1)) return 2;
 mt19937 rnd(120821011);
 while (true) {
  llu y = 2, yy = y, x = rnd() % n, t = 1;
 for (llu sz = 2; t == 1; sz <<= 1, y = yy) {
  for (llu i = 0; t == 1 && i < sz; ++i) {</pre>
    yy = f(yy, x, n);
    t = gcd(yy > y ? yy - y : y - yy, n);
  if (t != 1 && t != n) return t;
```

Polynomial Operations 5.19

```
using V = vector<int>;
#define fi(l, r) for (int i = int(l); i < int(r); ++i)
template <int mod, int G, int maxn> struct Poly : V {
 static uint32_t n2k(uint32_t n) {
 if (n <= 1) return 1;
  return 1u << (32 - __builtin_clz(n - 1));</pre>
 static NTT<mod,G,maxn> ntt; // coefficients in [0, P)
 explicit Poly(int n = 1) : V(n) {}
 Poly(const V &v) : V(v) {}
 Poly(const Poly &p, size_t n) : V(n) {
 copy_n(p.data(), min(p.size(), n), data());
 Poly &irev() { return reverse(data(), data() + size())
    , *this; }
 Poly &isz(int sz) { return resize(sz), *this; }
 Poly &iadd(const Poly &rhs) { // n() == rhs.n()
  fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
  return *this:
 Poly &imul(int k) {
  fi(0, size())(*this)[i] = modmul((*this)[i], k);
  return *this;
 Poly Mul(const Poly &rhs) const {
  const int sz = n2k(size() + rhs.size() - 1);
  Poly X(*this, sz), Y(rhs, sz);
  ntt(X.data(), sz), ntt(Y.data(), sz);
  fi(0, sz) X[i] = modmul(X[i], Y[i]);
  ntt(X.data(), sz, true);
  return X.isz(size() + rhs.size() - 1);
 Poly Inv() const { // coef[0] != 0
  if (size() == 1) return V{modinv(*begin())};
  const int sz = n2k(size() * 2);
  Poly X = Poly(*this, (size() + 1) / 2).Inv().isz(sz),
     Y(*this, sz);
  ntt(X.data(), sz), ntt(Y.data(), sz);
  fi(0, sz) X[i] = modmul(X[i], modsub(2, modmul(X[i],
    Y[i])));
  ntt(X.data(), sz, true);
  return X.isz(size());
 Poly Sqrt() const { // coef[0] \in [1, mod)^2
  if (size() == 1) return V{QuadraticResidue((*this))
    [0], mod)};
  Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
    size());
  return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
 pair<Poly, Poly> DivMod(const Poly &rhs) const {
  if (size() < rhs.size()) return {V{0}, *this};</pre>
  const int sz = size() - rhs.size() + 1;
  Poly X(rhs); X.irev().isz(sz);
  Poly Y(*this); Y.irev().isz(sz);
  Poly Q = Y.Mul(X.Inv()).isz(sz).irev();
  X = rhs.Mul(Q), Y = *this;
  fi(0, size()) Y[i] = modsub(Y[i], X[i]);
  return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
 Poly Dx() const {
```

```
Poly ret(size() - 1);
  fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
    1]);
  return ret.isz(max<int>(1, ret.size()));
Poly Sx() const {
 Poly ret(size() + 1);
 fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
    this)[i]);
 return ret:
Poly Ln() const { // coef[0] == 1
 return Dx().Mul(Inv()).Sx().isz(size());
                                                              }
Poly Exp() const { // coef[0] == 0
 if (size() == 1) return V{1};
 Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size)
    ()):
 Poly Y = X.Ln(); Y[0] = mod - 1;
  fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
 return X.Mul(Y).isz(size());
Poly Pow(const string &K) const {
 int nz = 0:
  while (nz < size() && !(*this)[nz]) ++nz;</pre>
  int nk = 0, nk2 = 0;
 for (char c : K) {
  nk = (nk * 10 + c - '0') \% mod;
  nk2 = nk2 * 10 + c - '0';
  if (nk2 * nz >= size())
   return Poly(size());
  nk2 %= mod - 1:
  if (!nk && !nk2) return Poly(V{1}, size());
 Poly X = V(data() + nz, data() + size() - nz * (nk2 - nz)
     1));
 int x0 = X[0];
  return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
    modpow(x0, nk2)).irev().isz(size()).irev();
V Eval(V x) const {
 if (x.empty()) return {};
 const size_t n = max(x.size(), size());
 vector<Poly> t(n * 2, V{1, 0}), f(n * 2);
for (size_t i = 0; i < x.size(); ++i)</pre>
   t[n + i] = V{1, mod-x[i]};
 for (size_t i = n - 1; i > 0; --i)
  t[i] = t[i * 2].Mul(t[i * 2 + 1]);
  f[1] = Poly(*this, n).irev().Mul(t[1].Inv()).isz(n).
    irev();
 for (size_t i = 1; i < n; ++i) {</pre>
   auto o = f[i]; auto sz = o.size();
  f[i*2] = o.irev().Mul(t[i*2+1]).isz(sz).irev().isz(t
    [i*2].size());
   f[i*2+1] = o.Mul(t[i*2]).isz(sz).irev().isz(t[i
    *2+1].size());
 for (size_t i=0;i<x.size();++i) x[i] = f[n+i][0];</pre>
 return x;
static int LinearRecursion(const V &a, const V &c,
    int64_t n) { // a_n = \sum_{i=1}^{n} a_i(n-j)
 const int k = (int)a.size();
  assert((int)c.size() == k + 1);
 Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
 fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
 while (n) {
  if (n % 2) W = W.Mul(M).DivMod(C).second;
  n /= 2, M = M.Mul(M).DivMod(C).second;
 int ret = 0;
 fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
 return ret;
}
#undef fi
using Poly_t = Poly<998244353, 3, 1 << 20>;
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
```

```
struct S {
 int MOD, w;
 int64_t x, y;
 S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
  : MOD(m), w(w_{-}), x(x_{-}), y(y_{-}) {}
 S operator*(const S &rhs) const {
  int w_ = w;
if (w_ == -1) w_ = rhs.w;
  assert(w_! = -1 \text{ and } w_! = rhs.w);
  return { MOD, w_,
   (x * rhs.x + y * rhs.y % MOD * w) % MOD,
   (x * rhs.y + y * rhs.x) % MOD };
int get_root(int n, int P) {
  if (P == 2 or n == 0) return n;
 auto check = [&](int x) {
  return qpow(x, (P - 1) / 2, P); };
if (check(n) != 1) return -1;
  int64_t a; int w; mt19937 rnd(7122);
  do { a = rnd() % P;
    w = ((a * a - n) \% P + P) \% P;
  } while (check(w) != P - 1);
  return qpow(S(P, w, a, 1), (P + 1) / 2).x;
5.21 Simplex
namespace simplex {
// maximize c^Tx under Ax <= B
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
 double inv = 1.0 / d[r][s];
 for (int i = 0; i < m + 2; ++i)
  for (int j = 0; j < n + 2; ++j)</pre>
   if (i != r && j != s)
    d[i][j] -= d[r][j] * d[i][s] * inv;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;</pre>
 for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
 d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
 int x = m + z;
 while (true) {
  int s = -1;
  for (int i = 0; i <= n; ++i) {</pre>
   if (!z && q[i] == -1) continue;
   if (s == -1 || d[x][i] < d[x][s]) s = i;</pre>
  if (d[x][s] > -eps) return true;
  int r = -1;
  for (int i = 0; i < m; ++i) {</pre>
   if (d[i][s] < eps) continue;</pre>
   if (r == -1 ||
    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
  pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
 m = b.size(), n = c.size();
 d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)</pre>
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)</pre>
  p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
 for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];</pre>
 q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
 for (int i = 1; i < m; ++i)</pre>
  if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
 if (d[r][n + 1] < -eps) {</pre>
```

pivot(r, n);

5.20 Quadratic residue

5.22 Simplex Construction

```
Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that for all 1 \leq j \leq m, \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j and x_i \geq 0 for all 1 \leq i \leq n.

1. In case of minimization, let c_i' = -c_i
2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \to \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
```

- $3. \sum_{1 \le i \le n} A_{ji} x_i = b_j$
 - · $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$ · $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6 Geometry

6.1 Basic Geometry

```
#define IM imag
#define RE real
using lld = int64_t;
using llf = long double;
using PT = std::complex<lld>;
using PTF = std::complex<llf>;
using P = PT;
auto toPTF(PT p) { return PTF{RE(p), IM(p)}; }
int sgn(lld x) \{ return (x > 0) - (x < 0); \}
lld dot(P a, P b) { return RE(conj(a) * b); }
lld cross(P a, P b) { return IM(conj(a) * b); }
int ori(P a, P b, P c) {
return sgn(cross(b - a, c - a));
namespace std {
bool operator<(const P &a, const P &b) {</pre>
return RE(a) != RE(b) ? RE(a) < RE(b) : IM(a) < IM(b);</pre>
int quad(P p) {
return (IM(p) == 0) // use sgn for PTF
  ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
int argCmp(P a, P b) {
// -1 / 0 / 1 <-> < / == / > (atan2)
int qa = quad(a), qb = quad(b);
if (qa != qb) return sgn(qa - qb);
return sgn(cross(b, a));
template <typename V> llf area(const V & pt) {
lld ret = 0;
for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
 ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
return ret / 2.0;
P rot90(P p) { return P{-IM(p), RE(p)}; }
PTF project(PTF p, PTF q) { // p onto q
return dot(p, q) * q / dot(q, q);
llf FMOD(llf x) {
if (x < -PI) x += PI * 2;
if (x > PI) x -= PI * 2;
return x;
```

6.2 Segment & Line Intersection

```
struct Segment { // closed segment
 PT st, dir; // represent st + t*dir for 0<=t<=1
 Segment(PT s, PT e) : st(s), dir(e - s) {}
 static bool valid(lld p, lld q) {
  // is there t s.t. 0 <= t <= 1 && qt == p ?
  if (q < 0) q = -q, p = -p;
  return 0 <= p && p <= q;
 vector<PT> ends() const { return { st, st + dir }; }
template <typename T> bool isInter(T A, PT P) {
 if (A.dir == PT(0)) return P == A.st; // BE CAREFUL
 return cross(P - A.st, A.dir) == 0 &&
  T::valid(dot(P - A.st, A.dir), norm(A.dir));
template <typename U, typename V>
bool isInter(U A, V B) {
   if (cross(A.dir, B.dir) == 0) { // BE CAREFUL
  bool res = false;
  for (PT P: A.ends()) res |= isInter(B, P);
  for (PT P: B.ends()) res |= isInter(A, P);
  return res;
 PT D = B.st - A.st;
lld C = cross(A.dir, B.dir);
 return U::valid(cross(D, B.dir), C) &&
  V::valid(cross(D, A.dir), C);
struct Line {
PT st, ed, dir;
Line (PT s, PT e)
  : st(s), ed(e), dir(e - s) {}
PTF intersect(const Line &A, const Line &B) {
llf t = cross(B.st - A.st, B.dir) /
  llf(cross(A.dir, B.dir));
 return toPTF(A.st) + PTF(t) * toPTF(A.dir);
```

6.3 2D Convex Hull

```
void make_hull(vector<pll> &dots) { // n=1 => ans = {}
  sort(dots.begin(), dots.end());
  vector<pll> ans(1, dots[0]);
  for (int ct = 0; ct < 2; ++ct, reverse(ALL(dots)))
    for (int i = 1, t = SZ(ans); i < SZ(dots); i++) {
      while (SZ(ans) > t && ori(
            ans[SZ(ans) - 2], ans.back(), dots[i]) <= 0)
            ans.pop_back();
      ans.pb(dots[i]);
    }
  ans.pop_back(), ans.swap(dots);
}</pre>
```

6.4 3D Convex Hull

```
// return the faces with pt indexes
int flag[MXN][MXN];
struct Point{
 ld x,y,z;
 Point operator * (const ld &b) const {
  return (Point) {x*b,y*b,z*b};}
 Point operator * (const Point &b) const {
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
};
Point ver(Point a, Point b, Point c) {
  return (b - a) * (c - a);}
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0;
 REP(i,n) REP(j,n) flag[i][j] = 0;
 vector<Face> now;
 now.emplace_back(0,1,2);
 now.emplace_back(2,1,0);
 for (int i=3; i<n; i++){</pre>
  ftop++; vector<Face> next;
  REP(j, SZ(now)) {
   Face& f=now[j]; int ff = 0;
   ld d=(pt[i]-pt[f.a]).dot(
     ver(pt[f.a], pt[f.b], pt[f.c]));
   if (d <= 0) next.push_back(f);</pre>
   if (d > 0) ff=ftop;
```

```
else if (d < 0) ff=-ftop;
  flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
}
REP(j, SZ(now)) {
  Face& f=now[j];
  if (flag[f.a][f.b] > 0 &&
     flag[f.a][f.b] != flag[f.b][f.a])
     next.emplace_back(f.a,f.b,i);
  if (flag[f.b][f.c] > 0 &&
     flag[f.b][f.c] != flag[f.c][f.b])
     next.emplace_back(f.b,f.c,i);
  if (flag[f.c][f.a] > 0 &&
     flag[f.c][f.a] != flag[f.a][f.c])
     next.emplace_back(f.c,f.a,i);
}
now=next;
}
return now;
}
```

6.5 2D Farthest Pair

```
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {
   while(abs(cross(stk[i+1]-stk[i],
        stk[(pos+1)%n]-stk[i])) >
        abs(cross(stk[i+1]-stk[i],
        stk[pos]-stk[i]))) pos = (pos+1)%n;
   ans = max({ans, dis(stk[i], stk[pos]),
        dis(stk[i+1], stk[pos]));
}
```

6.6 kD Closest Pair (3D ver.)

```
llf solve(vector<P> v) {
shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,</pre>
 unordered_map<lld, int>>> m;
llf d = dis(v[0], v[1]);
 auto Idx = [\&d] (llf x) \rightarrow lld {
 return round(x * 2 / d) + 0.1; };
auto rebuild_m = [&m, &v, &Idx](int k) {
 m.clear();
 for (int i = 0; i < k; ++i)
  m[Idx(v[i].x)][Idx(v[i].y)]
    [Idx(v[i].z)] = i;
}; rebuild_m(2);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
 const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
     kz = Idx(v[i].z); bool found = false;
  for (int dx = -2; dx <= 2; ++dx) {
   const lld nx = dx + kx;
   if (m.find(nx) == m.end()) continue;
   auto& mm = m[nx];
   for (int dy = -2; dy <= 2; ++dy) {
    const lld ny = dy + ky;
    if (mm.find(ny) == mm.end()) continue;
    auto& mmm = mm[ny];
    for (int dz = -2; dz \le 2; ++dz) {
     const lld nz = dz + kz;
     if (mmm.find(nz) == mmm.end()) continue;
     const int p = mmm[nz];
     if (dis(v[p], v[i]) < d) {</pre>
      d = dis(v[p], v[i]);
      found = true;
     }
  }
  if (found) rebuild_m(i + 1);
 else m[kx][ky][kz] = i;
return d;
```

6.7 Simulated Annealing

```
llf anneal() {
  mt19937 rnd_engine( seed );
  uniform_real_distribution< llf > rnd( 0, 1 );
  const llf dT = 0.001;
  // Argument p
```

```
llf S_cur = calc( p ), S_best = S_cur;
for ( llf T = 2000 ; T > EPS ; T -= dT ) {
    // Modify p to p_prime
    const llf S_prime = calc( p_prime );
    const llf delta_c = S_prime - S_cur;
    llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
    if ( rnd( rnd_engine ) <= prob )
        S_cur = S_prime, p = p_prime;
    if ( S_prime < S_best ) // find min
        S_best = S_prime, p_best = p_prime;
}
return S_best;
}</pre>
```

6.8 Half Plane Intersection

```
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
bool operator<(const Line &lhs, const Line &rhs) {</pre>
  if (int cmp = argCmp(lhs.dir, rhs.dir))
    return cmp == -1;
  return ori(lhs.st, lhs.ed, rhs.st) < 0;</pre>
// intersect function is in "Segment Intersect"
llf HPI(vector<Line> &lines) {
  sort(lines.begin(), lines.end());
  deque<Line> que;
  deque<PTF> pt;
  que.push_back(lines[0]);
  for (int i = 1; i < (int)lines.size(); i++) {</pre>
    if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
     continue;
#define POP(L, R) '
    while (pt.size() > 0 \
      && ori(L.st, L.ed, pt.back()) < 0) \
    pt.pop_back(), que.pop_back(); \
while (pt.size() > 0 \
      && ori(R.st, R.ed, pt.front()) < 0) \
      pt.pop_front(), que.pop_front();
    POP(lines[i], lines[i]);
    pt.push_back(intersect(que.back(), lines[i]));
    que.push_back(lines[i]);
  POP(que.front(), que.back())
  if (que.size() <= 1 ||
    argCmp(que.front().dir, que.back().dir) == 0)
    return 0;
  pt.push_back(intersect(que.front(), que.back()));
  return area(pt);
}
```

6.9 Minkowski Sum

6.10 Circle Class

```
struct Circle { PTF o; llf r; };
vector<llf> intersectAngle(Circle A, Circle B) {
   PTF dir = B.o - A.o; llf d2 = norm(dir);
   if (norm(A.r - B.r) >= d2) // norm(x) := |x|^2
   if (A.r < B.r) return {-PI, PI}; // A in B
   else return {}; // B in A
   if (norm(A.r + B.r) <= d2) return {};
   llf dis = abs(dir), theta = arg(dir);
   llf phi = acos((A.r * A.r + d2 - B.r * B.r) /
        (2 * A.r * dis));
   llf L = FMOD(theta - phi), R = FMOD(theta + phi);
   return { L, R };
   }
   vector<PTF> intersectPoint(Circle a, Circle b) {
```

```
llf d = abs(a.o - b.o);
if (d >= b.r+a.r || d <= abs(b.r-a.r)) return {};</pre>
llf dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
PTF dir = (a.o - b.o) / d;
PTF u = dir*d1 + b.o;
PTF v = rot90(dir) * sqrt(max<llf>(0, b.r*b.r-d1*d1));
return {u + v, u - v};
```

6.11 Intersection of line and Circle

```
vector<PTF> line_interCircle(const PTF &p1,
 const PTF &p2, const PTF &c, const double r) {
PTF ft = p1 + project(c-p1, p2-p1), vec = p2-p1;
llf dis = abs(c - ft);
if (abs(dis - r) < eps) return {ft};</pre>
if (dis > r) return {};
vec = vec * sqrt(r * r - dis * dis) / abs(vec);
return {ft + vec, ft - vec};
```

6.12 Intersection of Polygon and Circle

```
// Divides into multiple triangle, and sum up
  test by HDU2892
llf _area(PTF pa, PTF pb, llf r) {
if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
if (abs(pb) < eps) return 0;</pre>
llf S, h, theta;
llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
llf cosB = dot(pb, pb - pa) / a / c, B = acos(cosB);
llf cosC = dot(pa, pb) / a / b, C = acos(cosC);
if (a > r) {
 S = (C / 2) * r * r;
 h = a * b * sin(C) / c;
 if (h < r && B < PI / 2)
  S = (acos(h / r) * r * r - h * sqrt(r*r - h*h));
} else if (b > r) {
 theta = PI - B - asin(sin(B) / r * a);
 S = 0.5 * a*r * sin(theta) + (C - theta) / 2 * r*r;
} else
 S = 0.5 * sin(C) * a * b;
return S;
llf area_poly_circle(const vector<PTF> &poly,
 const PTF &0, const llf r) {
llf S = 0;
for (int i = 0, N = poly.size(); i < N; ++i)</pre>
 S += _area(poly[i] - 0, poly[(i + 1) % N] - 0, r) *
     ori(0, poly[i], poly[(i + 1) % N]);
return fabs(S);
```

6.13 Point & Hulls Tangent

```
#define above(P, Vi, Vj) (ori(P, Vi, Vj) > 0) // true
    if Vi is above Vi
#define below(P, Vi, Vj) (ori(P, Vi, Vj) < 0) // true</pre>
    if Vi is below Vj
// Rtangent_PointPolyC(): binary search for convex
    polygon right tangent
   Input: P = a 2D point (exterior to the polygon)
        n = number of polygon vertices
//
        V = array of vertices for a 2D convex polygon
    with V[n] = V[0]
   Return: index "i" of rightmost tangent point V[i]
int Rtangent_PointPolyC(PT P, int n, PT *V) {
if (n == 1) return 0;
if (below(P, V[1], V[0]) && !above(P, V[n - 1], V[0]))
for (int a = 0, b = n;;) {
 int c = (a + b) / 2, dnC = not above(P, V[c + 1], V[c
    1):
  if (dnC && !above(P, V[c - 1], V[c])) return c;
  if (above(P, V[a + 1], V[a]))
   ((dnC || above(P, V[a], V[c])) ? b : a) = c;
  else
   ((!dnC || !below(P, V[a], V[c])) ? a : b) = c;
}
// Ltangent_PointPolyC(): binary search for convex
    polygon left tangent
   Input: P = a \ 2D \ point \ (exterior \ to \ the \ polygon)
```

```
n = number of polygon vertices
        V = array of vertices for a 2D convex polygon
    with V[n]=V[0]
    Return: index "i" of leftmost tangent point V[i]
int Ltangent_PointPolyC(PT P, int n, PT *V) {
 if (n == 1) return 0;
 if (above(P, V[n - 1], V[0]) && !below(P, V[1], V[0]))
  return 0:
 for (int a = 0, b = n;;) {
  int c = (a + b) / 2, dnC = below(P, V[c + 1], V[c]);
  if (!below(P, V[c - 1], V[c]) && !dnC) return c;
  if (below(P, V[a + 1], V[a]))
   ((!dnC || below(P, V[a], V[c])) ? b : a) = c;
  else
   ((dnC | | !above(P, V[a], V[c])) ? a : b) = c;
```

6.14 Polygon Union

```
llf rat(P a, P b) { return sgn(RE(b)) ? llf(RE(a))/RE(b
    ) : llf(IM(a))/IM(b); }
llf polyUnion(vector<vector<P>>& poly) {
llf ret = 0; // area of poly[i] must be non-negative
 rep(i,0,sz(poly)) rep(v,0,sz(poly[i])) {
  P A = poly[i][v], B = poly[i][(v + 1) % sz(poly[i])];
  vector<pair<llf, int>> segs{{0, 0}, {1, 0}};
  rep(j,0,sz(poly)) if (i != j) {
   rep(u,0,sz(poly[j])) {
    P C = poly[j][u], D = poly[j][(u + 1) % sz(poly[j])
    if (int sc = ori(A, B, C), sd = ori(A, B, D); sc !=
     sd) {
     llf sa = cross(D-C, A-C), sb = cross(D-C, B-C);
     if (min(sc, sd) < 0)
      segs.emplace_back(sa / (sa - sb), sgn(sc - sd));
    } else if (!sc && !sd && j<i && sgn(dot(B-A,D-C))</pre>
     segs.emplace_back(rat(C - A, B - A), 1);
     segs.emplace_back(rat(D - A, B - A), -1);
  }
  sort(segs.begin(), segs.end());
  for (auto &s : segs) s.first = clamp<llf>(s.first, 0,
     1);;
  llf sum = 0;
  int cnt = segs[0].second;
  rep(j,1,sz(segs)) {
   if (!cnt) sum += segs[j].first - segs[j - 1].first;
   cnt += segs[j].second;
  ret += cross(A,B) * sum;
}
 return ret / 2;
```

6.15 Convex Hulls Tangent

```
// RLtangent_PolyPolyC(): get the RL tangent between
    two convex polygons
// Input:
   m = number of vertices in polygon 1
   V = array of vertices for convex polygon 1 with V[m
    n = number of vertices in polygon 2
    W = array of vertices for convex polygon 2 with W[n]
    ]=W[0]
// Output:
   *t1 = index of tangent point V[t1] for polygon 1
// *t2 = index of tangent point W[t2] for polygon 2
void RLtangent_PolyPolyC(int m, PT *V, int n, PT *W,
    int *t1, int *t2) {
 // * first get the initial vertex on each polygon
 // right tangent from W[0] to V
 int ix1 = Rtangent_PointPolyC(W[0], m, V);
 // left tangent from V[ix1] to W
 int ix2 = Ltangent_PointPolyC(V[ix1], n, W);
 // * ping-pong linear search until it stabilizes
 for (bool done = false; not done; ) {
  done = true; // assume done until...
  while (ori(W[ix2], V[ix1], V[ix1 + 1]) <= 0)</pre>
   ++ix1; // get Rtangent from W[ix2] to V
  while (ori(V[ix1], W[ix2], W[ix2 - 1]) >= 0) {
```

```
--ix2; // get Ltangent from V[ix1] to W
   done = false; // not done if had to adjust this
  }
 *t1 = ix1, *t2 = ix2;
}
```

6.16 Tangent line of Two Circle

```
tanline(const Circle &c1, const Circle &c2, int sign1){
// sign1 = 1 for outer tang, -1 for inter tang
 vector<Line> ret;
if (norm(c1.o - c2.o) < eps) return ret;</pre>
llf d = abs(c1.o - c2.o);
PTF v = (c2.o - c1.o) / d;
llf c = (c1.r - sign1 * c2.r) / d;
if (c * c > 1) return ret;
llf h = sqrt(max < llf > (0, 1 - c * c));
for (int sign2 : {1, -1}) {
 PTF n = c * v + sign2 * h * rot90(v);
 PTF p1 = c1.o + n * c1.r;
 PTF p2 = c2.o + n * (c2.r * sign1);
 if (norm(p2 - p1) < eps)
  p2 = p1 + rot90(c2.o - c1.o);
 ret.push_back({p1, p2});
return ret;
```

6.17 Minimum Covering Circle

```
template<typename P>
Circle getCircum(const P &a, const P &b, const P &c){
Real a\bar{1} = a.x-b.x, b1 = a.y-b.y;
Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
Real a2 = a.x-c.x, b2 = a.y-c.y;
Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
Circle cc;
cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
return cc;
}
template<typename P>
Circle MinCircleCover(const vector<P>& pts){
random_shuffle(pts.begin(), pts.end());
Circle c = { pts[0], 0 };
 for(int i=0;i<(int)pts.size();i++){</pre>
 if (dist(pts[i], c.o) <= c.r) continue;</pre>
  c = { pts[i], 0 };
 for (int j = 0; j < i; j++) {
  if(dist(pts[j], c.o) <= c.r) continue;
  c.o = (pts[i] + pts[j]) / 2;</pre>
   c.r = dist(pts[i], c.o);
   for (int k = 0; k < j; k++) {</pre>
    if (dist(pts[k], c.o) <= c.r) continue;</pre>
    c = getCircum(pts[i], pts[j], pts[k]);
   }
return c;
```

6.18 KDTree (Nearest Point)

```
const int MXN = 100005;
struct KDTree {
struct Node {
 int x,y,x1,y1,x2,y2;
 int id,f;
 Node *L, *R;
} tree[MXN], *root;
int n;
LL dis2(int x1, int y1, int x2, int y2) {
 LL dx = x1-x2, dy = y1-y2;
 return dx*dx+dy*dy;
static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
void init(vector<pair<int,int>> ip) {
 n = ip.size();
 for (int i=0; i<n; i++) {</pre>
```

```
tree[i].id = i;
   tree[i].x = ip[i].first;
   tree[i].y = ip[i].second;
  root = build_tree(0, n-1, 0);
 Node* build_tree(int L, int R, int d) {
  if (L>R) return nullptr;
  int M = (L+R)/2; tree[M].f = d%2;
  nth_element(tree+L,tree+M,tree+R+1,d%2?cmpy:cmpx);
  tree[M].x1 = tree[M].x2 = tree[M].x;
  tree[M].y1 = tree[M].y2 = tree[M].y;
  tree[M].L = build_tree(L, M-1, d+1);
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
   tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
  tree[M].R = build_tree(M+1, R, d+1);
  if (tree[M].R) {
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
   tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
   tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
  return tree+M;
 int touch(Node* r, int x, int y, LL d2){
  LL dis = sqrt(d2)+1;
  if (x<r->x1-dis || x>r->x2+dis ||
    y<r->y1-dis || y>r->y2+dis)
    return 0;
  return 1;
 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
  if (!r || !touch(r, x, y, md2)) return;
  LL d2 = dis2(r\rightarrow x, r\rightarrow y, x, y);
  if (d2 < md2 || (d2 == md2 \&\& mID < r->id)) {
   mID = r \rightarrow id;
   md2 = d2;
  // search order depends on split dim
  if ((r->f == 0 && x < r->x) ||
    (r->f == 1 \&\& y < r->y)) {
    nearest(r->L, x, y, mID, md2);
   nearest(r\rightarrow R, x, y, mID, md2);
  } else {
   nearest(r\rightarrow R, x, y, mID, md2);
   nearest(r\rightarrow L, x, y, mID, md2);
  }
 int query(int x, int y) {
  int id = 1029384756;
  LL d2 = 102938475612345678LL;
  nearest(root, x, y, id, d2);
  return id;
} tree;
        Rotating Sweep Line
```

```
void rotatingSweepLine(pair<int, int> a[], int n) {
 vector<pair<int, int>> l;
 l.reserve(n * (n - 1) / 2)
 for (int i = 0; i < n; ++i)</pre>
  for (int j = i + 1; j < n; ++j)
   l.emplace_back(i, j);
 sort(l.begin(), l.end(), [&a](auto &u, auto &v){
  lld udx = a[u.first].first - a[u.second].first;
  lld udy = a[u.first].second - a[u.second].second;
  lld vdx = a[v.first].first - a[v.second].first;
lld vdy = a[v.first].second - a[v.second].second;
  if (udx == 0 or vdx == 0) return not udx == 0;
  int s = sgn(udx * vdx);
  return udy * vdx * s < vdy * udx * s;</pre>
 });
 vector<int> idx(n), p(n);
 iota(idx.begin(), idx.end(), 0);
 sort(idx.begin(), idx.end(), [&a](int i, int j){
  return a[i] < a[j]; });</pre>
 for (int i = 0; i < n; ++i) p[idx[i]] = i;</pre>
 for (auto [i, j]: l) {
```

```
National Taiwan University - ckiseki
  // do here
 swap(p[i], p[j]);
 idx[p[i]] = i, idx[p[j]] = j;
6.20 Circle Cover
const int N = 1021;
struct CircleCover {
int C;
Cir c[N];
bool g[N][N], overlap[N][N];
// Area[i] : area covered by at least i circles
double Area[ N ];
void init(int _C){ C = _C;}
struct Teve {
 PTF p; double ang; int add;
 Teve() {}
 Teve(PTF _a, double _b, int _c):p(_a), ang(_b), add(
    c){}
 bool operator<(const Teve &a)const</pre>
  {return ang < a.ang;}
}eve[N * 2];
// strict: x = 0, otherwise x = -1
bool disjuct(Cir &a, Cir &b, int x)
 {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
bool contain(Cir &a, Cir &b, int x)
{return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
bool contain(int i, int j) {
 /* c[j] is non-strictly in c[i]. */
  return (sign(c[i].R - c[j].R) > 0 \mid | (sign(c[i].R - c
    [j].R) == 0 \&\& i < j)) \&\& contain(c[i], c[j], -1);
```

```
Stringology
7
```

7.1 Hash

}

}

void solve(){

++cnt;

else{

 $fill_n(Area, C + 2, 0);$

for(int i = 0; i < C; ++i)</pre>

for(int i = 0; i < C; ++i)</pre>

for(int j = 0; j < C; ++j)</pre>

for(int j = 0; j < C; ++j)</pre>

for(int i = 0; i < C; ++i){</pre>

for(int j = 0; j < C; ++j)</pre> if(i != j && g[i][j]) {

int E = 0, cnt = 1; for(int j = 0; j < C; ++j)</pre>

if(B > A) ++cnt;

sort(eve, eve + E);

cnt += eve[j].add;

eve[E] = eve[0];

overlap[i][j] = contain(i, j);

disjuct(c[i], c[j], -1));

if(j != i && overlap[j][i])

PTF aa = IP[0], bb = IP[1];

for(int j = 0; j < E; ++j){</pre>

if (theta < 0) theta += 2. * pi;</pre>

g[i][j] = !(overlap[i][j] || overlap[j][i] ||

auto IP = intersectPoint(c[i], c[j]);

llf A = arg(aa-c[i].0), B = arg(bb-c[i].0); eve[E++] = Teve(bb,B,1), eve[E++] = Teve(aa,A,-1);

if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;

Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;

Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;

double theta = eve[j + 1].ang - eve[j].ang;

```
class Hash {
private:
 static constexpr int P = 127, Q = 1051762951;
 vector<int> h, p;
public:
  void init(const string &s){
  h.assign(s.size()+1, 0); p.resize(s.size()+1);
```

```
for (size_t i = 0; i < s.size(); ++i)</pre>
    h[i + 1] = add(mul(h[i], P), s[i]);
   generate(p.begin(), p.end(),[x=1,y=1,this]()
     mutable{y=x;x=mul(x,P);return y;});
  int query(int l, int r){ // 1-base (l, r]
   return sub(h[r], mul(h[l], p[r-l]));}
};
```

7.2 Suffix Array

```
namespace sfx {
bool _t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], _c[maxn * 2];
int x[maxn], _p[maxn], _q[maxn * 2];
// sa[i]: sa[i]-th suffix is the
// i-th lexigraphically smallest suffix.
 // hi[i]: longest common prefix
 // of suffix sa[i] and suffix sa[i - 1].
void pre(int *a, int *c, int n, int z) {
 memset(a, 0, sizeof(int) * n);
 memcpy(x, c, sizeof(int) * z);
void induce(int *a,int *c,int *s,bool *t,int n,int z){
 memcpy(x + 1, c, sizeof(int) * (z - 1));
  for (int i = 0; i < n; ++i)
  if (a[i] && !t[a[i] - 1])
    a[x[s[a[i] - 1]]++] = a[i] - 1;
  memcpy(x, c, sizeof(int) * z);
  for (int i = n - 1; i >= 0; --i)
   if (a[i] && t[a[i] - 1])
    a[--x[s[a[i] - 1]]] = a[i] - 1;
 void sais(int *s, int *a, int *p, int *q,
 bool *t, int *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn=0, nmxz=-1, *nsa = a+n, *ns=s+n, last=-1;
  memset(c, 0, sizeof(int) * z);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];</pre>
  if (uniq) {
  for (int i = 0; i < n; ++i) a[--c[s[i]]] = i;</pre>
  return;
  for (int i = n - 2; i >= 0; --i)
  t[i] = (s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  pre(a, c, n, z);
for (int i = 1; i <= n - 1; ++i)</pre>
   if (t[i] && !t[i - 1])
    a[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(a, c, s, t, n, z);
for (int i = 0; i < n; ++i) {</pre>
   if (a[i] && t[a[i]] && !t[a[i] - 1]) {
   bool neq = last < 0 || \</pre>
    memcmp(s + a[i], s + last,
    (p[q[a[i]] + 1] - a[i]) * sizeof(int));
   ns[q[last = a[i]]] = nmxz += neq;
  }}
  sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
  pre(a, c, n, z);
  for (int i = nn - 1; i >= 0; --i)
  a[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(a, c, s, t, n, z);
void build(const string &s) {
  const int n = int(s.size());
  for (int i = 0; i < n; ++i) _s[i] = s[i];</pre>
  _s[n] = 0; // s shouldn't contain 0
 sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
for(int i = 0; i < n; ++i) rev[sa[i] = sa[i+1]] = i;</pre>
  int ind = hi[0] = 0;
  for (int i = 0; i < n; ++i) {</pre>
  if (!rev[i]) {
    ind = 0;
    continue;
  while (i + ind < n && \
    s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  hi[rev[i]] = ind ? ind-- : 0;
}}
```

7.3 Suffix Automaton

```
struct SuffixAutomaton {
struct node {
  int ch[K], len, fail, cnt, indeg;
 node(int L = 0) : ch{}, len(L), fail(0), cnt(0),
} st[N];
int root, last, tot;
void extend(int c) {
 int cur = ++tot;
  st[cur] = node(st[last].len + 1);
 while (last && !st[last].ch[c]) {
    st[last].ch[c] = cur;
    last = st[last].fail;
 if (!last) {
   st[cur].fail = root;
  } else {
    int q = st[last].ch[c];
    if (st[q].len == st[last].len + 1) {
      st[cur].fail = q;
    } else {
      int clone = ++tot;
      st[clone] = st[q];
      st[clone].len = st[last].len + 1;
      st[st[cur].fail = st[q].fail = clone].cnt = 0;
      while (last && st[last].ch[c] == q) {
        st[last].ch[c] = clone;
        last = st[last].fail;
      }
   }
 7
 st[last = cur].cnt += 1;
}
void init(const char* s) {
 root = last = tot = 1;
 st[root] = node(0);
  for (char c; c = *s; ++s) extend(c - 'a');
int q[N];
void dp() {
 for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
  int head = 0, tail = 0;
 for (int i = 1; i <= tot; i++)</pre>
    if (st[i].indeg == 0) q[tail++] = i;
 while (head != tail) {
    int now = q[head++];
    if (int f = st[now].fail) {
      st[f].cnt += st[now].cnt;
      if (--st[f].indeg == 0) q[tail++] = f;
   }
 }
int run(const char* s) {
 int now = root;
  for (char c; c = *s; ++s) {
   if (!st[now].ch[c -= 'a']) return 0;
   now = st[now].ch[c];
 return st[now].cnt;
}
} SAM;
```

7.4 Z value

```
vector<int> Zalgo(const string &s) {
  vector<int> z(s.size(), s.size());
  for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
    int j = clamp(r - i, 0, z[i - l]);
    for (; i + j < z[0] and s[i + j] == s[j]; ++j);
    if (i + (z[i] = j) > r) r = i + z[l = i];
  }
  return z;
}
```

7.5 Manacher

```
int z[maxn];
int manacher(const string& s) {
   string t = ".";
   for(char c: s) t += c, t += '.';
   int l = 0, r = 0, ans = 0;
```

```
for (int i = 1; i < t.length(); ++i) {
    z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
    while (i - z[i] >= 0 && i + z[i] < t.length()) {
        if(t[i - z[i]] == t[i + z[i]]) ++z[i];
        else break;
    }
    if (i + z[i] > r) r = i + z[i], l = i;
    }
    for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
    return ans;
}</pre>
```

7.6 Lexico Smallest Rotation

```
string mcp(string s) {
  int n = s.length();
  s += s; int i = 0, j = 1;
  while (i < n && j < n) {
    int k = 0;
    while (k < n && s[i + k] == s[j + k]) k++;
    ((s[i + k] <= s[j + k]) ? j : i) += k + 1;
    j += (i == j);
  }
  return s.substr(i < n ? i : j, n);
}</pre>
```

7.7 Main Lorentz

```
vector<tuple<tuple<size_t, size_t, int, int>>> reps;
void find_repetitions(const string &s, int shift = 0) {
 if (s.size() <= 1)
  return;
 const size_t nu = s.size() / 2, nv = s.size() - nu;
 string u = s.substr(0, nu), v = s.substr(nu);
 string ru(u.rbegin(), u.rend());
 string rv(v.rbegin(), v.rend());
 find_repetitions(u, shift);
 find_repetitions(v, shift + nu);
 auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u),
z3 = Zalgo(ru + '#' + rv), z4 = Zalgo(v);
 for (size_t cntr = 0; cntr < s.size(); cntr++) {</pre>
  size_t l; int k1, k2;
  if (cntr < nu) {</pre>
   l = nu - cntr;
   k1 = l < z1.size() ? z1[l] : 0;
   k2 = n + 1 - l < z2.size() ? z2[n + 1 - l] : 0;
  } else {
   l = cntr - nu + 1;
   k1 = n + 1 - l < z3.size() ? z3[n + 1 - l] : 0;
   k2 = l < z4.size() ? z4[l] : 0;
  if (k1 + k2 >= 1)
   reps.emplace_back(cntr, l, k1, k2);
}
```

7.8 BWT

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a'
 vector<int> v[ SIGMA ];
 void BWT(char* ori, char* res){
  // make ori -> ori + ori
    then build suffix array
 void iBWT(char* ori, char* res){
  for( int i = 0 ; i < SIGMA ; i ++ )</pre>
   v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )</pre>
   v[ ori[i] - BASE ].push_back( i );
  vector<int> a;
  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
   for( auto j : v[ i ] ){
    a.push_back( j );
    ori[ ptr ++ ] = BASE + i;
  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){</pre>
   res[ i ] = ori[ a[ ptr ] ];
   ptr = a[ ptr ];
  res[ len ] = 0;
} bwt;
```

7.9 Palindromic Tree

```
struct palindromic_tree{
struct node{
  int next[26],f,len;
  int cnt,num,st,ed; // num = depth of fail link
  node(int l=0):f(0),len(l),cnt(0),num(0) {
  memset(next, 0, sizeof(next)); }
};
vector<node> st;
vector<char> s;
int last,n;
void init(){
 st.clear();s.clear();last=1; n=0;
  st.push_back(0);st.push_back(-1);
  st[0].f=1;s.push_back(-1); }
int getFail(int x){
 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
  return x;}
 void add(int c){
 s.push_back(c='a'); ++n;
  int cur=getFail(last);
  if(!st[cur].next[c]){
   int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
   st[now].num=st[st[now].f].num+1;
  last=st[cur].next[c];
  ++st[last].cnt;}
 void dpcnt() { // cnt = #occurence in whole str
  for (int i=st.size()-1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
int size(){ return st.size()-2;}
} pt;
int main() {
string s; cin >> s; pt.init();
for (int i=0; i<SZ(s); i++) {</pre>
  int prvsz = pt.size(); pt.add(s[i]);
  if (prvsz != pt.size()) {
   int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [l,r]: s.substr(l, r-l+1)
 }
return 0;
```

8 Misc

8.1 Theorems

8.1.1 Sherman-Morrison formula

$$(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1+v^{\mathsf{T}}A^{-1}u}$$

8.1.2 Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i), L_{ij} = -c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

8.1.3 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij}=x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij}=-d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.4 Cayley's Formula

- Given a degree sequence d_1,d_2,\ldots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.5 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if

 $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \le k \le n$.

8.1.6 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.7 Euler's planar graph formula

 $V - E + F = C + 1, E \le 3V - 6$ (?)

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2)$, find maximum $S\in I_1\cap I_2$. For each iteration, build the directed graph and find a shortest path from s to t.

```
\begin{array}{l} \cdot \ s \to x : S \sqcup \{x\} \in I_1 \\ \cdot \ x \to t : S \sqcup \{x\} \in I_2 \\ \\ \cdot \ y \to x : S \setminus \{y\} \sqcup \{x\} \in I_1 \ (y \ \text{is in the unique circuit of} \ S \sqcup \{x\}) \\ \cdot \ x \to y : S \setminus \{y\} \sqcup \{x\} \in I_2 \ (y \ \text{is in the unique circuit of} \ S \sqcup \{x\}) \end{array}
```

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|.$ In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S,$ resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

8.2 Bitset LCS

```
scanf("%d%d", &n, &m), u = n / 64 + 1;
for (int i = 1, c; i <= n; i++)
  scanf("%d", &c), p[c].set(i);
for (int i = 1, c; i <= m; i++) {
  scanf("%d", &c), (g = f) |= p[c];
  f.shiftLeftByOne(), f.set(0);
  ((f = g - f) ^= g) &= g;
}
printf("%d\n", f.count());</pre>
```

8.3 Prefix Substring LCS

```
void all_lcs(string s, string t) { // 0-base
vector<int> h(SZ(t));
iota(ALL(h), 0);
for (int a = 0; a < SZ(s); ++a) {
  int v = -1;
  for (int c = 0; c < SZ(t); ++c)
   if (s[a] == t[c] || h[c] < v)
      swap(h[c], v);
  // LCS(s[0, a], t[b, c]) =
  // c - b + 1 - sum([h[i] >= b] | i <= c)
  // h[i] might become -1 !!
}</pre>
```

8.4 Convex 1D/1D DP

```
struct segment {
 int i, l, r;
 segment() {}
 segment(int a, int b, int c): i(a), l(b), r(c) {}
inline lld f(int l, int r){return dp[l] + w(l+1, r);}
void solve() {
 dp[0] = 0;
 deque<segment> dq; dq.push_back(segment(0, 1, n));
 for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i);
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
   f(i, dq.back().l)<f(dq.back().i, dq.back().l))</pre>
    dq.pop_back();
  if (dq.size()) {
   int d = 1 << 20, c = dq.back().l;
while (d >>= 1) if (c + d <= dq.back().r)</pre>
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
```

```
dq.back().r = c; seg.l = c + 1;
}
if (seg.l <= n) dq.push_back(seg);
}
</pre>
```

8.5 ConvexHull Optimization

```
struct L {
 mutable int64_t a, b, p;
 bool operator<(const L &r) const {</pre>
  return a < r.a; /* here */ }</pre>
 bool operator<(int64_t x) const { return p < x; }</pre>
struct DynamicHull : multiset<L, less<>>> {
 static const int64_t kInf = 1e18;
 bool Isect(iterator x, iterator y) {
  auto Div = [](int64_t a, int64_t b) {
    return a / b - ((a ^ b) < 0 && a % b); };
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a)
   x\rightarrow p = x\rightarrow b > y\rightarrow b ? kInf : -kInf; /* here */
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(int64_t a, int64_t b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
  while ((y = x) != begin() && (--x)->p >= y->p)
   Isect(x, erase(y));
 int64_t Query(int64_t x) { // default chmax
  auto l = *lower_bound(x); // to chmin:
  return l.a * x + l.b; // modify the 2 "<>"
};
```

8.6 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

8.7 Tree Knapsack

```
int dp[N][K]; PII obj[N];
vector<int> G[N];
void dfs(int u, int mx){
    for(int s: G[u]) {
        if(mx < obj[s].first) continue;
        for(int i=0;i<=mx-obj[s].FF;i++)
        dp[s][i] = dp[u][i];
        dfs(s, mx - obj[s].first);
        for(int i=obj[s].FF;i<=mx;i++)
        dp[u][i] = max(dp[u][i],
        dp[s][i - obj[s].FF] + obj[s].SS);
    }
}</pre>
```

8.8 N Queens Problem

```
vector< int > solve( int n ) {
   // no solution when n=2, 3
   vector< int > ret;
   if ( n % 6 == 2 ) {
      for ( int i = 2 ; i <= n ; i += 2 )
        ret.push_back( i );
   ret.push_back( 3 ); ret.push_back( 1 );
   for ( int i = 7 ; i <= n ; i += 2 )
      ret.push_back( i );
   ret.push_back( i );
   ret.push_back( 5 );
   } else if ( n % 6 == 3 ) {
      for ( int i = 4 ; i <= n ; i += 2 )
      ret.push_back( i );
   ret.push_back( i );
```

```
ret.push_back( 2 );
for ( int i = 5 ; i <= n ; i += 2 )
    ret.push_back( i );
ret.push_back( 1 ); ret.push_back( 3 );
} else {
    for ( int i = 2 ; i <= n ; i += 2 )
        ret.push_back( i );
    for ( int i = 1 ; i <= n ; i += 2 )
        ret.push_back( i );
}
return ret;
}</pre>
```

8.9 Stable Marriage

```
1: Initialize m \in M and w \in W to free
2: while \exists free man m who has a woman w to propose to do
        w \leftarrow \text{first woman on } m's list to whom m has not yet proposed
4:
       if \exists some pair (m', w) then
           if w prefers m to m' then
6:
              m' \leftarrow free
               (m, w) \leftarrow \mathsf{engaged}
           end if
       else
10:
            (m, w) \leftarrow \mathsf{engaged}
11:
       end if
12: end while
```

8.10 Binary Search On Fraction

```
struct Q {
 ll p, q;
 Q go(Q b, ll d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
 Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
 assert(pred(hi));
 bool dir = 1, L = 1, H = 1;
 for (; L || H; dir = !dir) {
  ll len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
   if (Q mid = hi.go(lo, len + step);
     \label{eq:mid.p} \mbox{mid.p} > \mbox{N} \ || \ \mbox{mid.q} > \mbox{N} \ || \ \mbox{dir} \ \ ^{\mbox{pred(mid))}}
    t++;
   else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
 return dir ? hi : lo;
```