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| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6  |  | 1 1 1 1          | 3 3 3 3 4 4 4                                   |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7   |  | 1 1 1 1 1 1 1 1  | 3 3 3 3 4 4 4 4                                 |
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| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9   |  | 1                | 3 3 3 3 4 4 4 4 4                               |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10   |  | 1                | 3 3 3 3 3 4 4 4 4 4 5                           |
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| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12   | L <sup>n</sup> / <sub>3</sub>   Enumeration           Strling Number           5.2.1 First Kind           5.2.2 Second Kind           ax+by=gcd           Berlekamp Massey           Charateristic Polynomial           Chinese Remainder           De-Bruijn           DiscreteLog           Extended Euler           ExtendedFloorSum           Fast Fourier Transform           FloorSum  | 1                | 33333444445555                                  |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12   | Lm₁/3   Enumeration           Strling Number           5.2.1 First Kind           5.2.2 Second Kind           ax+by=gcd           Berlekamp Massey           Charateristic Polynomial           Chinese Remainder           De-Bruijn           DiscreteLog           Extended Euler           ExtendedFloorSum           Fast Fourier Transform   | 1                | 3 3 3 3 3 4 4 4 4 4 5 5                         |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13   | L <sup>n</sup> / <sub>3</sub>   Enumeration           Strling Number           5.2.1 First Kind           5.2.2 Second Kind           ax+by=gcd           Berlekamp Massey           Charateristic Polynomial           Chinese Remainder           De-Bruijn           DiscreteLog           Extended Euler           ExtendedFloorSum           Fast Fourier Transform           FloorSum  | 1                | 33333444445555                                  |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13<br>5.14   | L <sup>n</sup> / <sub>3</sub>   Enumeration           Strling Number           5.2.1 First Kind           5.2.2 Second Kind           ax+by=gcd           Berlekamp Massey           Charateristic Polynomial           Chinese Remainder           De-Bruijn           DiscreteLog           Extended Euler           ExtendedFloorSum           Fast Fourier Transform           FloorSum           FWT  | 1                | 333334444445556                                 |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.12<br>5.12<br>5.13<br>5.14<br>5.15   | Lm₁/3   Enumeration           Strling Number           5.2.1 First Kind           5.2.2 Second Kind           αx+by=gcd           Berlekamp Massey           Charateristic Polynomial           Chinese Remainder           De-Bruijn           DiscreteLog           Extended Euler           ExtendedFloorSum           Fast Fourier Transform           FloorSum           FWT           Miller Rabin           NTT   | 1                | 33333444445556666                               |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16   | Lm₁/3   Enumeration           Strling Number           5.2.1 First Kind           5.2.2 Second Kind           xx+by=gcd           Berlekamp Massey           Charateristic Polynomial           Chinese Remainder           De-Bruijn           DiscreteLog           Extended Euler           ExtendedFloorSum           Fast Fourier Transform           FloorSum           FWT           Miller Rabin           NTT           Partition Number  | 1                | 333334444455566666                              |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17   | L <sup>n</sup> / <sub>3</sub>   Enumeration           Strling Number           5.2.1 First Kind           5.2.2 Second Kind           ax+by=gcd           Berlekamp Massey           Charateristic Polynomial           Chinese Remainder           De-Bruijn           DiscreteLog           Extended Euler           Extended FloorSum           Fast Fourier Transform           FloorSum           FWT           Miller Rabin           NTT           Partition Number           Pi Count (Linear Sieve)   | 1                | 3333344444555666666                             |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18   | L <sup>n</sup> / <sub>3</sub>   Enumeration           Strling Number           5.2.1 First Kind           5.2.2 Second Kind           ax+by=gcd           Berlekamp Massey           Charateristic Polynomial           Chinese Remainder           De-Bruijn           DiscreteLog           Extended Euler           ExtendedFloorSum           Fast Fourier Transform           FloorSum           FWT           Miller Rabin           NTT           Partition Number           Pi Count (Linear Sieve)           Pollard Rho  | 1                | 3 3 3 3 3 4 4 4 4 4 4 5 5 5 6 6 6 6 6 7         |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19   | L n/2   Enumeration           Strling Number           5.2.1 First Kind           5.2.2 Second Kind           αx+by=gcd           Berlekamp Massey           Charateristic Polynomial           Chinese Remainder           De-Bruijn           DiscreteLog           Extended Euler           ExtendedFloorSum           Fast Fourier Transform           FloorSum           FWT           Miller Rabin           NTT           Partition Number           Pi Count (Linear Sieve)           Pollard Rho           Polynomial Operations  | 1                | 3 3 3 3 3 4 4 4 4 4 4 5 5 5 6 6 6 6 6 7 7       |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20   | Lm/3   Enumeration         Strling Number         5.2.1 First Kind         5.2.2 Second Kind         αx+by=gcd         Berlekamp Massey         Charateristic Polynomial         Chinese Remainder         De-Bruijn         DiscreteLog         Extended Euler         ExtendedFloorSum         Fast Fourier Transform         FloorSum         FWT         Miller Rabin         NTT         Partition Number         Pi Count (Linear Sieve)         Pollard Rho         Polynomial Operations         O Quadratic residue   | 1                | 3 3 3 3 3 4 4 4 4 4 4 5 5 5 6 6 6 6 6 7 7 7     |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.10<br>5.11<br>5.12<br>5.13   | Lm₁/3   Enumeration           Strling Number           5.2.1 First Kind           5.2.2 Second Kind           xx+by=gcd           Berlekamp Massey           Charateristic Polynomial           Chinese Remainder           De-Bruijn           DiscreteLog           Extended Euler           ExtendedFloorSum           Fast Fourier Transform           FloorSum           FWT           Miller Rabin           NTT           Partition Number           Pi Count (Linear Sieve)           Pollard Rho           Polynomial Operations           Quadratic residue           Simplex  | 1                | 3 3 3 3 3 4 4 4 4 4 4 5 5 5 6 6 6 6 6 7 7 7 8   |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.10<br>5.11<br>5.12<br>5.13   | Lm/3   Enumeration         Strling Number         5.2.1 First Kind         5.2.2 Second Kind         αx+by=gcd         Berlekamp Massey         Charateristic Polynomial         Chinese Remainder         De-Bruijn         DiscreteLog         Extended Euler         ExtendedFloorSum         Fast Fourier Transform         FloorSum         FWT         Miller Rabin         NTT         Partition Number         Pi Count (Linear Sieve)         Pollard Rho         Polynomial Operations         O Quadratic residue   | 1                | 3 3 3 3 3 4 4 4 4 4 4 5 5 5 6 6 6 6 6 7 7 7     |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>5.20   | Lm/3   Enumeration           Strling Number           5.2.1 First Kind           5.2.2 Second Kind           αx+by=gcd           Berlekamp Massey           Charateristic Polynomial           Chinese Remainder           De-Bruijn           DiscreteLog           Extended Euler           ExtendedFloorSum           Fast Fourier Transform           FloorSum           FWT           Miller Rabin           NTT           Partition Number           Pi Count (Linear Sieve)           Pollard Rho           Polynomial Operations           O Quadratic residue           Simplex           2 Simplex Construction  | 1                | 3 3 3 3 3 4 4 4 4 4 4 5 5 5 6 6 6 6 6 7 7 7 8 8 |
| 5  | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>5.21<br>5.21<br>5.21<br>5.21<br>5.21<br>5.21<br>5.21<br>5.21   | L note   | 1                | 33333444445556666677788                         |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>5.21<br>5.22<br>Geol. 6.1  | L m/3   Enumeration         Strling Number         5.2.1 First Kind         5.2.2 Second Kind         αx+by=gcd         Berlekamp Massey         Charateristic Polynomial         Chinese Remainder         De-Bruijn         DiscreteLog         Extended Euler         ExtendedFloorSum         Fast Fourier Transform         FloorSum         FWT         Miller Rabin         NTT         Partition Number         Pi Count (Linear Sieve)         Pollard Rho         Polynomial Operations         Quadratic residue         Simplex         2 Simplex Construction         Image: Provided Sieve Polynomial Construction         Image: Provided Sieve Polyn   | 1                | 33333444445556666677788 88                      |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.12<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>5.21<br>5.21<br>5.21<br>5.21<br>5.21<br>5.21<br>5.21<br>5.21   | Lm₁/3   Enumeration         Strling Number         5.2.1 First Kind         5.2.2 Second Kind         ax+by=gcd         Berlekamp Massey         Charateristic Polynomial         Chinese Remainder         De-Bruijn         DiscreteLog         Extended Euler         Extended FloorSum         Fast Fourier Transform         FloorSum         FWT         Miller Rabin         NTT         Partition Number         Pi Count (Linear Sieve)         Pollard Rho         Polynomial Operations         Quadratic residue         Simplex         2 Simplex Construction         Ometry         Basic Geometry         Segment & Line Intersection  | 1                | 33333444445556666677788                         |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>5.21<br>5.22<br>Geol. 6.1  | L Table 1   Enumeration   Strling Number   5.2.1 First Kind   5.2.2 Second Kind   3x+by=gcd   Berlekamp Massey   Charateristic Polynomial   Chinese Remainder   De-Bruijn   DiscreteLog   Extended Euler   Extended FloorSum   Fast Fourier Transform   FloorSum   FWT   Miller Rabin   NTT   Partition Number   Pi Count (Linear Sieve)   Pollard Rho   Polynomial Operations   O Quadratic residue   Simplex   Simpl | 1                | 33333444445556666677788 88                      |
|    | 5.1<br>5.2<br>5.3<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>5.21<br>5.22<br>Geol  | Lm₁/3   Enumeration         Strling Number         5.2.1 First Kind         5.2.2 Second Kind         ax+by=gcd         Berlekamp Massey         Charateristic Polynomial         Chinese Remainder         De-Bruijn         DiscreteLog         Extended Euler         Extended FloorSum         Fast Fourier Transform         FloorSum         FWT         Miller Rabin         NTT         Partition Number         Pi Count (Linear Sieve)         Pollard Rho         Polynomial Operations         Quadratic residue         Simplex         2 Simplex Construction         Ometry         Basic Geometry         Segment & Line Intersection  | 1                | 333334444455566666677788 8888                   |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>5.21<br>5.22<br>6.21<br>6.2<br>6.2<br>6.3   | L Table 1   Enumeration   Strling Number   5.2.1 First Kind   5.2.2 Second Kind   3x+by=gcd   Berlekamp Massey   Charateristic Polynomial   Chinese Remainder   De-Bruijn   DiscreteLog   Extended Euler   Extended FloorSum   Fast Fourier Transform   FloorSum   FWT   Miller Rabin   NTT   Partition Number   Pi Count (Linear Sieve)   Pollard Rho   Polynomial Operations   O Quadratic residue   Simplex   Simpl | 1                | 33333444445556666677788 88889                   |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4   | L m/3   Enumeration         Strling Number         5.2.1 First Kind         5.2.2 Second Kind         αx+by=gcd         Berlekamp Massey         Charateristic Polynomial         Chinese Remainder         De-Bruijn         DiscreteLog         Extended Euler         Extended FloorSum         Fast Fourier Transform         FloorSum         FWT         Miller Rabin         NTT         Partition Number         Pi Count (Linear Sieve)         Pollard Rho         Polynomial Operations         Quadratic residue         Simplex         2 Simplex Construction         Demetry         Basic Geometry         Segment & Line Intersection         2D Convex Hull         3D Convex Hull         2D Farthest Pair  | 1                | 333334444445556666677788 88899                  |
|    | 5.1<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.10<br>5.11<br>5.12<br>5.14<br>5.15<br>5.16<br>5.17<br>5.19<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.5<br>6.6  | L m/3   Enumeration         Strling Number         5.2.1 First Kind         5.2.2 Second Kind         αx+by=gcd         Berlekamp Massey         Charateristic Polynomial         Chinese Remainder         De-Bruijn         DiscreteLog         Extended Euler         Extended FloorSum         Fast Fourier Transform         FloorSum         FWT         Miller Rabin         NTT         Partition Number         Pi Count (Linear Sieve)         Pollard Rho         Polynomial Operations         Quadratic residue         Simplex         2 Simplex Construction         Immetry         Basic Geometry         Segment & Line Intersection         2D Convex Hull         3D Convex Hull         2D Farthest Pair         kD Closest Pair (3D ver.)  | 1                | 33333444445556666677788 8889999                 |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.19<br>5.10<br>5.11<br>5.12<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.5<br>6.5<br>6.7  | L m/3  | 1                | 33333444445556666677788 88899999                |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.6<br>6.7<br>6.8   | L Tips Number 5.2.1 First Kind 5.2.2 Second Kind 5.2.2 Second Kind 5.2.2 Second Kind 5.2.3 Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations O Quadratic residue Simplex Simplex Construction Demetry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 3D Convex Hull 2D Farthest Pair KD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection  | 1                | 333334444455566666677788 888999990              |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.5<br>6.6<br>6.7<br>6.8<br>6.9  | L m/3   Enumeration         Strling Number         5.2.1 First Kind         5.2.2 Second Kind         ax+by=gcd         Berlekamp Massey         Charateristic Polynomial         Chinese Remainder         De-Bruijn         DiscreteLog         Extended Euler         Extended FloorSum         Fast Fourier Transform         FloorSum         FWT         Miller Rabin         NTT         Partition Number         Pi Count (Linear Sieve)         Pollard Rho         Polynomial Operations         Quadratic residue         Simplex         Simplex Construction         Ometry         Basic Geometry         Segment & Line Intersection         2D Convex Hull         3D Convex Hull         3D Convex Hull         2D Farthest Pair         kD Closest Pair (3D ver.)         Simulated Annealing         Half Plane Intersection         Minkowski Sum  | 1                | 33333444445556666677788 8889999900              |
|    | 5.1<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.5<br>6.6<br>6.7<br>6.9<br>6.9<br>6.9<br>6.9<br>6.9<br>6.9<br>6.9<br>6.9<br>6.9<br>6.9   | L m/3   Enumeration         Strling Number         5.2.1 First Kind         5.2.2 Second Kind         αx+by=gcd         Berlekamp Massey         Charateristic Polynomial         Chinese Remainder         De-Bruijn         DiscreteLog         Extended Euler         Extended FloorSum         Fast Fourier Transform         FloorSum         FWT         Miller Rabin         NTT         Partition Number         Pi Count (Linear Sieve)         Pollard Rho         Polynomial Operations         Quadratic residue         Simplex         2 Simplex Construction         Demetry         Basic Geometry         Segment & Line Intersection         2D Convex Hull         2D Farthest Pair         kD Closest Pair (3D ver.)         Simulated Annealing         Half Plane Intersection         Minkowski Sum         Circle Class  | 1                | 33333444445556666677788 88899999000             |
|    | 5.1<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.10<br>5.11<br>5.12<br>5.14<br>5.15<br>5.16<br>5.17<br>5.19<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.5<br>6.6<br>6.7<br>6.8<br>6.9<br>6.0<br>6.0<br>6.0<br>6.0<br>6.0<br>6.0<br>6.0<br>6.0<br>6.0<br>6.0   | Lm₁ Strling Number         5.2.1 First Kind         5.2.2 Second Kind         αx+by=gcd         Berlekamp Massey         Charateristic Polynomial         Chinese Remainder         De-Bruijn         DiscreteLog         Extended Euler         ExtendedFloorSum         Fast Fourier Transform         FloorSum         FWT         Miller Rabin         NTT         Partition Number         Pi Count (Linear Sieve)         Pollard Rho         Polynomial Operations         Quadratic residue         Simplex         2 Simplex Construction         Immetry         Basic Geometry         Segment & Line Intersection         2D Convex Hull         3D Convex Hull         3D Convex Hull         3D Convex Hull         2D Farthest Pair         kD Closest Pair (3D ver.)         Simulated Annealing         Half Plane Intersection         Minkowski Sum         Circle Class         Intersection of line and Circle  | 1                | 333334444455566666677788 888999990000           |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.10<br>5.11<br>5.12<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.6<br>6.7<br>6.8<br>6.9<br>6.10<br>6.11<br>6.12<br>6.12<br>6.13<br>6.14<br>6.14<br>6.15<br>6.15<br>6.15<br>6.15<br>6.15<br>6.15<br>6.15<br>6.15 | Time      | 1                | 333334444455566666677788 888999990000           |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.10<br>5.11<br>5.12<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.6<br>6.7<br>6.8<br>6.9<br>6.10<br>6.11<br>6.12<br>6.12<br>6.13<br>6.14<br>6.14<br>6.15<br>6.15<br>6.15<br>6.15<br>6.15<br>6.15<br>6.15<br>6.15 | Lm₁ Strling Number         5.2.1 First Kind         5.2.2 Second Kind         αx+by=gcd         Berlekamp Massey         Charateristic Polynomial         Chinese Remainder         De-Bruijn         DiscreteLog         Extended Euler         ExtendedFloorSum         Fast Fourier Transform         FloorSum         FWT         Miller Rabin         NTT         Partition Number         Pi Count (Linear Sieve)         Pollard Rho         Polynomial Operations         Quadratic residue         Simplex         2 Simplex Construction         Immetry         Basic Geometry         Segment & Line Intersection         2D Convex Hull         3D Convex Hull         3D Convex Hull         3D Convex Hull         2D Farthest Pair         kD Closest Pair (3D ver.)         Simulated Annealing         Half Plane Intersection         Minkowski Sum         Circle Class         Intersection of line and Circle  | 1                | 33333444445556666677788 8889999900000           |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.9<br>5.10<br>5.11<br>5.12<br>5.13<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.5<br>6.6<br>6.7<br>6.8<br>6.9<br>6.10<br>6.10<br>6.10<br>6.10<br>6.10<br>6.10<br>6.10<br>6.10                                   | Time      | 1                | 33333444445556666677788 8889999900000           |
|    | 5.1<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.8<br>5.9<br>5.10<br>5.11<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.5<br>6.6<br>6.1<br>6.11<br>6.12<br>6.13<br>6.14<br>6.14<br>6.14<br>6.15<br>6.15<br>6.15<br>6.16<br>6.16<br>6.16<br>6.16<br>6.16                 | Partition Number   Partition   | 1                | 33333444445556666677788 88899999000000          |
|    | 5.1<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.9<br>5.10<br>5.11<br>5.12<br>5.14<br>5.15<br>5.16<br>5.17<br>5.18<br>5.19<br>5.20<br>6.1<br>6.2<br>6.4<br>6.6<br>6.7<br>6.6<br>6.7<br>6.1<br>6.1<br>6.1<br>6.1<br>6.1<br>6.1<br>6.1<br>6.1<br>6.1<br>6.1  | Part   Enumeration   Strling Number   5.2.1 First Kind   5.2.2 Second Kind   αx+by=gcd   Berlekamp Massey   Charateristic Polynomial   Chinese Remainder   De-Bruijn   DiscreteLog   Extended Euler   ExtendedFloorSum   Fost Fourier Transform   FloorSum   FWT   Miller Rabin   NTT   Partition Number   Pi Count (Linear Sieve)   Pollard Rho   Polynomial Operations   Quadratic residue   Simplex   Simplex   Simplex Construction   Description   Descript   | 1                | 33333444445556666677788 8889999900000011        |
|    | 5.1<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.10<br>5.11<br>5.12<br>5.14<br>5.15<br>5.16<br>5.17<br>5.19<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.6<br>6.7<br>6.8<br>6.10<br>6.11<br>6.12<br>6.13<br>6.14<br>6.15<br>6.15<br>6.15<br>6.16<br>6.16<br>6.16<br>6.16<br>6.16                               | Partial   Enumeration   Strling Number   S.2.1   First Kind   S.2.2   Second Kind   S.   | 1                | 333334444455566666677788 888999990000001111     |
|    | 5.1<br>5.2<br>5.3<br>5.4<br>5.5<br>5.6<br>5.7<br>5.19<br>5.12<br>5.14<br>5.15<br>5.16<br>5.17<br>5.19<br>5.20<br>6.1<br>6.2<br>6.3<br>6.4<br>6.6<br>6.7<br>6.8<br>6.9<br>6.10<br>6.11<br>6.12<br>6.13<br>6.14<br>6.14<br>6.15<br>6.16<br>6.16<br>6.16<br>6.16<br>6.16<br>6.16<br>6.17<br>6.17         | Part   Enumeration   Strling Number   5.2.1 First Kind   5.2.2 Second Kind   αx+by=gcd   Berlekamp Massey   Charateristic Polynomial   Chinese Remainder   De-Bruijn   DiscreteLog   Extended Euler   ExtendedFloorSum   Fost Fourier Transform   FloorSum   FWT   Miller Rabin   NTT   Partition Number   Pi Count (Linear Sieve)   Pollard Rho   Polynomial Operations   Quadratic residue   Simplex   Simplex   Simplex Construction   Description   Descript   | 1                | 333334444455566666677788 88899999000000111111   |

```
1
7 Stringology
 7.1
   22
 23
 23
  BWT .....
 8 Misc
 8.1 Theorems
   8.1.1 Kirchhoff's Theorem
8.1.2 Tutte's Matrix
     8.1.3

      8.1.5
      Havel-Hakimi algorithm

      8.1.6
      Euler's planar graph formula

   8.1.7
     1
  Basic
1.1 vimrc
se is nu ru et tgc sc hls cin cino+=j1 sw=4 sts=4 bs=2
  mouse=a encoding=utf-8 ls=2
colorscheme desert
filetype indent on
inoremap {<CR> {<CR>}<ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -std=c++17 -
  DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
  Wconversion -fsanitize=address,undefined -g && echo
   success<CR>
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 &&
   echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
1.2 Debug Macro
```

```
#ifdef KISEKI
#define safe cerr<<__PRETTY_FUNCTION__\</pre>
  <<" line "<<__LINE__<<" safe\n"
#define debug(a...) qwerty(#a, a)
#define orange(a...) dvorak(#a, a)
using std::cerr;
template <typename ...T>
void qwerty(const char *s, T ...a) {
  cerr << "\e[1;32m(" << s << ") = (";
  int cnt = sizeof...(T);
  (..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
template <typename Iter>
void dvorak(const char *s, Iter L, Iter R) {
  cerr << "\e[1;32m[ " << s << " ] = [ ";</pre>
  for (int f = 0; L != R; ++L)
    cerr << (f++?", " : "") << *L;
    cerr << " ]\e[0m\n";
#else
#define safe ((void)0)
#define debug(...) ((void)\theta)
#define orange(...) ((void)0)
#endif
```

### 1.3 Increase Stack

```
const int size = 256 << 20;</pre>
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));
```

# 1.4 Pragma Optimization

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8000)
```

```
1.5 IO Optimization
```

```
static inline int gc() {
  constexpr int B = 1<<20;
  static char buf[B], *p, *q;
  if(p == q &&
    (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
    return EOF;
  return *p++;
}

template < typename T >
  static inline bool gn( T &x ) {
  int c = gc(); T sgn = 1; x = 0;
  while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
  if(c == '-') sgn = -1, c = gc();
  if(c == EOF) return false;
  while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
  return x *= sgn, true;
}</pre>
```

# 2 Data Structure

# 2.1 Dark Magic

```
2.2 Link-Cut Tree
template <typename Val> class LCT {
private:
struct node
 int pa, ch[2];
 bool rev;
 Val v, v_prod, v_rprod;
 node() : pa{0}, ch{0, 0}, rev{false}, v{}, v_prod{},
    v_rprod{} {};
vector<node> nodes;
 set<pair<int, int>> edges;
bool is_root(int u) const {
 const int p = nodes[u].pa;
  return nodes[p].ch[0] != u and nodes[p].ch[1] != u;
bool is_rch(int u) const {
 return (not is_root(u)) and nodes[nodes[u].pa].ch[1]
    == u;
void down(int u) {
 if (auto &cnode = nodes[u]; cnode.rev) {
  if (cnode.ch[0]) set_rev(cnode.ch[0]);
   if (cnode.ch[1]) set_rev(cnode.ch[1]);
   cnode.rev = false;
}
void up(int u) {
 auto &cnode = nodes[u];
 cnode.v_prod =
  nodes[cnode.ch[0]].v_prod * cnode.v * nodes[cnode.ch
    [1]].v_prod;
 cnode.v_rprod =
  nodes[cnode.ch[1]].v_rprod * cnode.v * nodes[cnode.
    ch[0]].v_rprod;
void set_rev(int u) {
 swap(nodes[u].ch[0],\ nodes[u].ch[1]);\\
  swap(nodes[u].v_prod, nodes[u].v_rprod);
 nodes[u].rev ^= 1;
 void rotate(int u) {
 int f = nodes[u].pa, g = nodes[f].pa, l = is_rch(u);
```

```
if (nodes[u].ch[1 ^ 1])
   nodes[nodes[u].ch[1 ^ 1]].pa = f;
  if (not is_root(f))
   nodes[g].ch[is_rch(f)] = u;
  nodes[f].ch[1] = nodes[u].ch[1 ^ 1];
  nodes[u].ch[1 ^ 1] = f
  nodes[u].pa = g, nodes[f].pa = u;
  up(f);
 void splay(int u) {
  vector<int> stk = {u};
  while (not is_root(stk.back()))
   stk.push_back(nodes[stk.back()].pa);
  for (; not stk.empty(); stk.pop_back())
   down(stk.back());
  for(int f=nodes[u].pa;!is_root(u);f=nodes[u].pa){
   if(!is_root(f))rotate(is_rch(u)==is_rch(f)?f:u);
   rotate(u):
  }
  up(u);
 void access(int u) {
  int last = 0;
  for (int last = 0; u; last = u, u = nodes[u].pa) {
   splay(u);
   nodes[u].ch[1] = last;
   up(u);
 int find_root(int u) {
  access(u); splay(u);
  int la = 0:
  for (; u; la = u, u = nodes[u].ch[0]) down(u);
  return la;
 void change_root(int u) {
  access(u); splay(u); set_rev(u);
 void link(int x, int y)
  change_root(y); nodes[y].pa = x;
 void split(int x, int y) {
  change_root(x); access(y); splay(y);
 void cut(int x, int y) {
  split(x, y)
  nodes[y].ch[0] = nodes[x].pa = 0;
  up(y);
public:
 LCT(int n = 0) : nodes(n + 1) {}
 int add(const Val &v = {}) {
 nodes.push_back(v);
  return int(nodes.size()) - 2;
 int add(Val &&v) {
  nodes.emplace_back(move(v));
  return int(nodes.size()) - 2;
 void set_val(int u, const Val &v) {
  splay(++u); nodes[u].v = v; up(u);
 Val query(int x, int y) {
  split(++x, ++y);
  return nodes[y].v_prod;
 bool connected(int u, int v) { return find_root(++u)
    == find_root(++v);
 void add_edge(int u, int v) {
  if (++u > ++v) swap(u, v)
  edges.emplace(u, v); link(u, v);
 void del_edge(int u, int v) {
  auto k = minmax(++u, ++v)
  if (auto it = edges.find(k); it != edges.end()) {
   edges.erase(it); cut(u, v);
  }
};
```

# .3 LiChao Segment Tree

struct L {

```
int m, k, id;
                                                                #undef sz
 L(): id(-1) {}
 L(int a, int b, int c) : m(a), k(b), id(c) {}
                                                               2.5 Linear Basis
 int at(int x) { return m * x + k; }
                                                               template <int BITS>
class LiChao {
                                                               struct LinearBasis {
private:
                                                                array<uint64_t, BITS> basis;
 int n; vector<L> nodes;
                                                                Basis() { basis.fill(0); }
                                                                void add(uint64_t x)
 static int lc(int x) { return 2 * x + 1; }
 static int rc(int x) { return 2 * x + 2;
                                                                 for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
  if (basis[i] == 0) {
 void insert(int 1, int r, int id, L ln) {
  int m = (1 + r) >> 1;
                                                                   basis[i] = x;
  if (nodes[id].id == -1) {
                                                                    return;
   nodes[id] = ln;
   return;
                                                                  x ^= basis[i];
                                                                 }
  bool atLeft = nodes[id].at(1) < ln.at(1);</pre>
  if (nodes[id].at(m) < ln.at(m)) {</pre>
                                                                bool ok(uint64_t x) {
   atLeft ^= 1;
                                                                 for (int i = 0; i < BITS; ++i)</pre>
   swap(nodes[id], ln);
                                                                  if ((x >> i) & 1) x ^= basis[i];
                                                                 return x == 0;
  if (r - 1 == 1) return;
  if (atLeft) insert(1, m, lc(id), ln);
                                                               };
  else insert(m, r, rc(id), ln);
                                                               2.6
                                                                      Binary Search On Segment Tree
 int query(int 1, int r, int id, int x) {
  int ret = 0, m = (1 + r) >> 1;
                                                               // find_first = x -> minimal x s.t. check( [a, x) )
                                                               // find_last = x \rightarrow maximal x s.t. check([x, b))
  if (nodes[id].id != -1)
                                                               template <typename C>
   ret = nodes[id].at(x);
                                                               int find_first(int 1, const C &check) {
  if (r - 1 == 1) return ret;
                                                                if (1 >= n) return n + 1;
  if (x < m) return max(ret, query(1, m, lc(id), x));</pre>
                                                                1 += sz;
  return max(ret, query(m, r, rc(id), x));
                                                                for (int i = height; i > 0; i--)
                                                                 propagate(1 >> i);
                                                                Monoid sum = identity;
public:
 LiChao(int n_{-}) : n(n_{-}), nodes(n * 4) {}
                                                                 while ((1 & 1) == 0) 1 >>= 1;
 void insert(L ln) { insert(0, n, 0, ln); }
                                                                 if (check(f(sum, data[1]))) {
 int query(int x) { return query(0, n, 0, x); }
                                                                  while (1 < sz) {</pre>
                                                                   propagate(1);
                                                                    1 <<= 1;
2.4 Treap
                                                                    auto nxt = f(sum, data[1]);
namespace Treap{
                                                                    if (not check(nxt)) {
 #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                                     sum = nxt;
 struct node{
                                                                     1++;
  int size;
                                                                   }
  uint32_t pri;
                                                                  }
  node *lc, *rc, *pa;
                                                                   return 1 + 1 - sz;
  node():size(0),pri(rand()),lc(0),rc(0),pa(0){}
  void pull() {
                                                                 sum = f(sum, data[1++]);
  size = 1; pa = nullptr;
                                                                } while ((1 & -1) != 1);
   if ( lc ) { size += lc->size; lc->pa = this; }
if ( rc ) { size += rc->size; rc->pa = this; }
                                                                return n + 1;
  }
                                                               template <typename C>
                                                               int find_last(int r, const C &check) {
node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
                                                                if (r <= 0) return -1;
                                                                r += sz;
  if ( L->pri > R->pri ) {
                                                                for (int i = height; i > 0; i--)
  L->rc = merge( L->rc, R ); L->pull();
                                                                 propagate((r - 1) >> i);
   return L;
                                                                Monoid sum = identity;
  } else {
                                                                do {
   R->lc = merge( L, R->lc ); R->pull();
   return R;
                                                                 while (r > 1 \text{ and } (r \& 1)) r >>= 1;
  }
                                                                 if (check(f(data[r], sum))) {
 }
                                                                  while (r < sz) {
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                                   propagate(r);
  if ( not rt ) L = R = nullptr;
                                                                    r = (r << 1) + 1;
  else if( sz( rt->lc ) + 1 <= k ) {
                                                                    auto nxt = f(data[r], sum);
                                                                    if (not check(nxt)) {
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                     sum = nxt;
   L->pull();
                                                                     r--:
  } else {
                                                                   }
   R = rt:
   split_by_size( rt->lc, k, L, R->lc );
                                                                  return r - sz;
   R->pull();
                                                                 }
                                                                sum = f(data[r], sum);
} while ((r & -r) != r);
 } // sz(L) == k
 int getRank(node *o) { // 1-base
int r = sz(o->lc) + 1;
                                                                return -1;
  for (;o->pa != nullptr; o = o->pa)
   if (o->pa->rc == o) r += sz(o->pa->lc) + 1;
                                                                    Graph
  return r;
```

3.1 2-SAT (SCC)

G.clear(); G.resize(n = n\_);

```
class TwoSat{
                                                                 low.assign(n, ecnt = 0);
private:
                                                                dfn.assign(n, 0);
 int n;
                                                               }
 vector<vector<int>> rG,G,sccs;
                                                               void add_edge(int u, int v) {
                                                                G[u].emplace_back(v, ecnt);
 vector<int> ord,idx;
 vector<bool> vis,result;
                                                                G[v].emplace_back(u, ecnt++);
 void dfs(int u){
                                                               void solve() {
  vis[u]=true
  for(int v:G[u])
                                                                bridge.assign(ecnt, false);
                                                                for (int i = 0; i < n; ++i)
   if(!vis[v]) dfs(v);
                                                                 if (not dfn[i]) dfs(i, i);
  ord.push_back(u);
 void rdfs(int u){
                                                               bool is_bridge(int x) { return bridge[x]; }
                                                             } bcc_bridge;
 vis[u]=false;idx[u]=sccs.size()-1;
  sccs.back().push_back(u);
                                                             3.3 BCC Vertex
  for(int v:rG[u])
   if(vis[v])rdfs(v);
                                                             class BCC_AP {
                                                              private:
public:
                                                               int n, ecnt;
 void init(int n_){
                                                               vector<vector<pair<int,int>>> G;
                                                               vector<int> bcc, dfn, low, st;
 G.clear();G.resize(n=n_);
                                                               vector<bool> ap, ins;
void dfs(int u, int f)
  rG.clear();rG.resize(n);
  sccs.clear();ord.clear();
                                                                dfn[u] = low[u] = dfn[f] + 1;
  idx.resize(n);result.resize(n);
                                                                 int ch = 0;
                                                                for (auto [v, t]: G[u]) if (v != f) {
  if (not ins[t]) {
 void add_edge(int u,int v){
 G[u].push_back(v);rG[v].push_back(u);
                                                                  st.push_back(t);
                                                                   ins[t] = true;
 void orr(int x,int y){
  if ((x^y)==1)return
  add_edge(x^1,y); add_edge(y^1,x);
                                                                  if (dfn[v]) {
                                                                  low[u] = min(low[u], dfn[v]);
 bool solve(){
                                                                   continue:
  vis.clear();vis.resize(n);
                                                                  } ++ch; dfs(v, u);
  for(int i=0;i<n;++i)</pre>
                                                                  low[u] = min(low[u], low[v]);
   if(not vis[i])dfs(i);
                                                                  if (low[v] >= dfn[u]) {
                                                                  ap[u] = true;
  reverse(ord.begin(),ord.end());
  for (int u:ord){
                                                                   while (true) {
   if(!vis[u])continue;
                                                                    int eid = st.back(); st.pop_back();
   sccs.push_back(vector<int>());
                                                                    bcc[eid] = ecnt;
   rdfs(u);
                                                                    if (eid == t) break;
  for(int i=0;i<n;i+=2)</pre>
                                                                  ecnt++;
   if(idx[i]==idx[i+1])
                                                                 }
    return false;
  vector<bool> c(sccs.size());
                                                                 if (ch == 1 and u == f) ap[u] = false;
  for(size_t i=0;i<sccs.size();++i){</pre>
                                                              public:
   for(auto sij : sccs[i]){
    result[sij]=c[i];
                                                               void init(int n_) {
    c[idx[sij^1]]=!c[i];
                                                                G.clear(); G.resize(n = n_);
   }
                                                                 ecnt = 0; ap.assign(n, false);
                                                                low.assign(n, 0); dfn.assign(n, 0);
  return true;
                                                               void add_edge(int u, int v) {
                                                                G[u].emplace_back(v, ecnt);
G[v].emplace_back(u, ecnt++);
 bool get(int x){return result[x];}
 int get_id(int x){return idx[x];}
 int count(){return sccs.size();}
} sat2;
                                                               void solve() {
                                                                ins.assign(ecnt, false);
      BCC Edge
3.2
                                                                bcc.resize(ecnt); ecnt = 0;
                                                                for (int i = 0; i < n; ++i)
if (not dfn[i]) dfs(i, i);</pre>
class BCC_Bridge {
 private:
  int n, ecnt;
                                                               int get_id(int x) { return bcc[x]; }
  vector<vector<pair<int,int>>> G;
  vector<int> dfn, low;
                                                               int count() { return ecnt;
  vector<bool> bridge;
                                                               bool is_ap(int x) { return ap[x]; }
  void dfs(int u, int f)
                                                             } bcc_ap;
   dfn[u] = low[u] = dfn[f] + 1;
                                                             3.4 Centroid Decomposition
   for (auto [v, t]: G[u]) {
    if (v == f) continue;
                                                             struct Centroid {
    if (dfn[v]) {
                                                              vector<vector<int64_t>> Dist;
                                                              vector<int> Parent, Depth;
     low[u] = min(low[u], dfn[v]);
     continue;
                                                              vector<int64_t> Sub, Sub2;
                                                              vector<int> Sz, Sz2;
    dfs(v, u);
                                                              Centroid(vector<vector<pair<int, int>>> g) {
    low[u] = min(low[u], low[v]);
                                                               int N = g.size();
                                                               vector<bool> Vis(N);
    if (low[v] > dfn[u]) bridge[t] = true;
                                                               vector<int> sz(N), mx(N);
                                                               vector<int> Path;
                                                               Dist.resize(N)
 public:
  void init(int n_) {
                                                               Parent.resize(N);
```

Depth.resize(N);

in[e.v] = e.w;

```
auto DfsSz = [&](auto dfs, int x) -> void {
                                                                   prv[e.v] = e.u;
   Vis[x] = true; sz[x] = 1; mx[x] = 0;
   for (auto [u, w] : g[x]) {
                                                                 in[root] = 0;
    if (Vis[u]) continue;
                                                                 prv[root] = -1;
    dfs(dfs, u)
                                                                 for (int i = 0; i < n; i++)
                                                                  if (in[i] == -inf)
    sz[x] += sz[u];
    mx[x] = max(mx[x], sz[u]);
                                                                   return -inf;
                                                                  // find cycle
   Path.push_back(x);
                                                                  int tot = 0;
                                                                 vector<int> id(n, -1), vis(n, -1);
for (int i = 0; i < n; i++) {</pre>
  }:
  auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
                                                                  ans += in[i];
   Dist[x].push_back(D);Vis[x] = true;
                                                                  for (int x = i; x != -1 && id[x] == -1; x = prv[x])
   for (auto [u, w] : g[x]) {
   if (Vis[u]) continue;
                                                                   if (vis[x] == i) {
    dfs(dfs, u, D + w);
                                                                     for (int y = prv[x]; y != x; y = prv[y])
                                                                      id[y] = tot;
  };
                                                                     id[x] = tot++;
  auto Dfs = [&]
                                                                    break;
   (auto dfs, int x, int D = 0, int p = -1)->void {
   Path.clear(); DfsSz(DfsSz, x);
                                                                   vis[x] = i;
   int M = Path.size();
                                                                  }
   int C = -1;
   for (int u : Path) {
                                                                 if (!tot)
                                                                  return ans;
    if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
    Vis[u] = false;
                                                                 for (int i = 0; i < n; i++)</pre>
                                                                  if (id[i] == -1)
   DfsDist(DfsDist, C);
                                                                   id[i] = tot++;
   for (int u : Path) Vis[u] = false;
                                                                  // shrink
   Parent[C] = p; Vis[C] = true;
                                                                 for (auto &e : E) {
   Depth[C] = D;
                                                                  if (id[e.u] != id[e.v])
   for (auto [u, w] : g[C]) {
                                                                   e.w -= in[e.v];
    if (Vis[u]) continue
                                                                  e.u = id[e.u], e.v = id[e.v];
    dfs(dfs, u, D + 1, C);
                                                                 n = tot;
                                                                 root = id[root];
  Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
  Sz.resize(N); Sz2.resize(N);
                                                                assert(false);
void Mark(int v) {
                                                              } DMST:
  int x = v, z = -1
                                                              3.6 Dominator Tree
 for (int i = Depth[v]; i >= 0; --i) {
Sub[x] += Dist[v][i]; Sz[x]++;
                                                              namespace dominator {
  if (z != -1) {
                                                              vector<int> g[maxn], r[maxn], rdom[maxn];
                                                              int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
    Sub2[z] += Dist[v][i];
    Sz2[z]++;
                                                              void init(int n) {
   z = x; x = Parent[x];
                                                               // vertices are numbered from 0 to n-1
  }
                                                               fill(dfn, dfn + n, -1);fill(rev, rev + n, -1);
                                                               fill(fa, fa + n, -1); fill(val, val + n, -1);
int64_t Query(int v) {
                                                               fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
                                                               fill(dom, dom + n, -1); tk = 0;
for (int i = 0; i < n; ++i) {
 int64_t res = 0;
 int x = v, z = -1;
 for (int i = Depth[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
                                                                g[i].clear(); r[i].clear(); rdom[i].clear();
  if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
  z = x; x = Parent[x];
                                                              void add_edge(int x, int y) { g[x].push_back(y); }
                                                              void dfs(int x) {
                                                               rev[dfn[x] = tk] = x;
  return res;
                                                               fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
                                                               for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
};
3.5 Directed Minimum Spanning Tree
                                                                r[dfn[u]].push_back(dfn[x]);
struct DirectedMST { // find maximum
                                                               }
struct Edge {
                                                              void merge(int x, int y) { fa[x] = y; }
 int u, v;
                                                              int find(int x, int c = 0) {
  int w;
                                                               if (fa[x] == x) return c ? -1 : x;
 Edge(int u, int v, int w) : u(u), v(v), w(w) {}
                                                               int p = find(fa[x], 1);
                                                               if (p == -1) return c ? fa[x] : val[x];
vector<Edge> Edges;
                                                               if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
void clear() { Edges.clear(); }
void addEdge(int a, int b, int w) { Edges.emplace_back
                                                               fa[x] = p
                                                               return c ? p : val[x];
    (a, b, w); }
int solve(int root, int n) {
  vector<Edge> E = Edges;
                                                              vector<int> build(int s, int n) {
                                                              // return the father of each node in the dominator tree
  int ans = 0:
 while (true) {
                                                              // p[i] = -2 if i is unreachable from s
   // find best in edge
                                                               dfs(s);
   vector<int> in(n, -inf), prv(n, -1);
                                                               for (int i = tk - 1; i >= 0; --i) {
                                                                for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
                                                                if (i) rdom[sdom[i]].push_back(i);
```

for (int &u : rdom[i]) {

chain[u] = chain[v];

```
int p = find(u);
   if (sdom[p] == i) dom[u] = i;
                                                                    if (chain[u] == 0) chain[u] = ++chains;
   else dom[u] = p;
                                                                   void dfschain(int u, int f) {
                                                                    tl[u] = timer++
 if (i) merge(i, rp[i]);
                                                                    if (head[chain[u]] == -1)
vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)</pre>
                                                                     head[chain[u]] = u;
                                                                    for (int v : G[u])
  if (v != f and chain[v] == chain[u])
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
                                                                      dfschain(v, u);
                                                                    for (int v : G[u])
 return p;
                                                                     if (v != f and chain[v] != chain[u])
                                                                      dfschain(v, u);
3.7 Edge Coloring
                                                                    tr[u] = timer;
// max(d_u) + 1 edge coloring, time: O(NM)
                                                                 public:
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
                                                                   LBD(int n) : timer(0), chains(0), G(n), tl(n), tr(n),
void clear(int N) {
                                                                      chain(n), \ head(n, \ -1), \ dep(n), \ pa(n) \ \{\}
for (int i = 0; i <= N; i++)
 for (int j = 0; j <= N; j++)
C[i][j] = G[i][j] = 0;
                                                                   void add_edge(int u, int v) {
                                                                    G[u].push_back(v); G[v].push_back(u);
void solve(vector<pair<int, int>> &E, int N) {
                                                                   void decompose() { predfs(0, 0); dfschain(0, 0); }
int X[kN] = {}, a;
auto update = [&](int u) {
                                                                   PII get_subtree(int u) { return {tl[u], tr[u]}; }
                                                                   vector<PII> get_path(int u, int v) {
                                                                    vector<PII> res;
while (chain[u] != chain[v]) {
 for (X[u] = 1; C[u][X[u]]; X[u]++);
auto color = [&](int u, int v, int c) {
                                                                     if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
 int p = G[u][v];
                                                                      swap(u, v)
                                                                     int s = head[chain[u]];
res.emplace_back(tl[s], tl[u] + 1);
 G[u][v] = G[v][u] = c;
 C[u][c] = v, C[v][c] = u;
 C[u][p] = C[v][p] = 0;
                                                                     u = pa[s];
 if (p) X[u] = X[v] = p;
  else update(u), update(v);
                                                                    if (dep[u] < dep[v]) swap(u, v);</pre>
                                                                    res.emplace_back(tl[v], tl[u] + 1);
  return p;
                                                                    return res;
auto flip = [&](int u, int c1, int c2) {
 int p = C[u][c1];
                                                                 };
 swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
if (!C[u][c1]) X[u] = c1;
                                                                        Manhattan Minimum Spanning Tree
                                                                 typedef Point<int> P:
  if (!C[u][c2]) X[u] = c2;
                                                                 vector<array<int, 3>> manhattanMST(vector<P> ps) {
                                                                   vi id(sz(ps));
 return p;
                                                                   iota(all(id), 0);
for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {
  auto [u, v] = E[t];</pre>
                                                                   vector<array<int, 3>> edges;
                                                                   rep(k, 0, 4) {
  sort(all(id),
                                                                                   [&](int i, int j) {
  int v0 = v, c = X[u], c0 = c, d;
                                                                     return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
 vector<pair<int, int>> L; int vst[kN] = {};
                                                                    });
  while (!G[u][v0]) {
                                                                    map<int, int> sweep;
   L.emplace_back(v, d = X[v]);
if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
                                                                    for (int i : id) {
                                                                     for (auto it = sweep.lower_bound(-ps[i].y);
     c = color(u, L[a].first, c);
                                                                        it != sweep.end(); sweep.erase(it++)) {
                                                                      int j = it->second
   else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
     color(u, L[a].first, L[a].second);
                                                                      P d = ps[i] - ps[j];
   else if (vst[d]) break
                                                                      if (d.y > d.x) break;
   else vst[d] = 1, v = C[u][d];
                                                                      edges.push_back({d.y + d.x, i, j});
 if (!G[u][v0]) {
  for (; v; v = flip(v, c, d), swap(c, d));
                                                                     sweep[-ps[i].y] = i;
   if (C[u][c0]) { a = int(L.size()) - 1;
                                                                    for (P &p : ps)
    while (--a >= 0 && L[a].second != c);
                                                                     if (k \& 1) p.x = -p.x;
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
                                                                     else swap(p.x, p.y);
   } else t--;
                                                                   return edges; // [{w, i, j}, ...]
                                                                 }
                                                                 3.10 MaxClique
3.8 Lowbit Decomposition
                                                                 // contain a self loop u to u, than u won't in clique
                                                                 template < size_t MAXN >
class LBD {
int timer, chains;
                                                                 class MaxClique{
vector<vector<int>> G;
                                                                 private:
vector<int> t1, tr, chain, head, dep, pa;
                                                                   using bits = bitset< MAXN >;
 // chains : number of chain
                                                                   bits popped, G[ MAXN ], ans;
// tl, tr[u] : subtree interval in the seq. of u
                                                                   size_t deg[ MAXN ], deo[ MAXN ], n;
 // head[i] : head of the chain i
                                                                   void sort_by_degree() {
 // chian[u] : chain id of the chain u is on
                                                                    popped.reset();
void predfs(int u, int f) {
  dep[u] = dep[pa[u] = f] + 1;
                                                                    for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                      deg[ i ] = G[ i ].count();
                                                                    for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
  for (int v : G[u]) if (v != f) {
                                                                      size_t mi = MAXN, id = 0;
   predfs(v, u);
                                                                      for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )</pre>
   if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
```

p < kN; p = int(cs[k].\_Find\_next(p))) {

```
mi = deg[ id = j ];
                                                                    r[t] = p; c[t++] = k;
    popped[ deo[ i ] = id ] = 1;
    for( size_t u = G[ i ]._Find_first() ;
     u < n ; u = G[ i ]._Find_next( u ) )</pre>
      -- deg[ u ];
                                                                 void dfs(vector<int> &r, vector<int> &c, int 1,
  }
                                                                  bitset<kN> mask) {
                                                                  while (!r.empty()) {
 void BK( bits R, bits P, bits X ) {
                                                                   int p = r.back(); r.pop_back();
  if (R.count()+P.count() <= ans.count()) return;</pre>
                                                                   mask[p] = 0;
  if ( not P.count() and not X.count() ) {
                                                                   if (q + c.back() <= ans) return;</pre>
  if ( R.count() > ans.count() ) ans = R;
                                                                   cur[q++] = p;
                                                                   vector<int> nr, nc;
   return;
                                                                   bitset<kN> nmask = mask & a[p];
  }
  /* greedily chosse max degree as pivot
                                                                   for (int i : r)
  bits cur = P | X; size_t pivot = 0, sz = 0;
                                                                    if (a[p][i]) nr.push_back(i);
  for ( size_t u = cur._Find_first() ;
                                                                   if (!nr.empty()) {
   u < n ; u = cur._Find_next( u )</pre>
                                                                    if (1 < 4) {
   if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
                                                                     for (int i : nr)
  cur = P & ( ~G[ pivot ] );
                                                                       d[i] = int((a[i] & nmask).count());
  */ // or simply choose first
                                                                      sort(nr.begin(), nr.end(),
  bits cur = P & (~G[ ( P | X )._Find_first() ]);
                                                                       [&](int x, int y)
  for ( size_t u = cur._Find_first()
                                                                        return d[x] > d[y];
   u < n ; u = cur._Find_next( u ) ) {
if ( R[ u ] ) continue;</pre>
                                                                       });
                                                                   csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
   R[u] = 1;
   BK(R, P & G[u], X & G[u]);
R[u] = P[u] = 0, X[u] = 1;
                                                                    ans = q; copy(cur, cur + q, sol);
                                                                   c.pop_back(); q--;
public:
                                                                  }
 void init( size_t n_ ) {
                                                                 int solve(bitset<kN> mask) { // vertex mask
  n = n_{-};
                                                                  vector<int> r, c;
for (int i = 0; i < n; i++)</pre>
  for ( size_t i = 0 ; i < n ; ++ i )
   G[ i ].reset();
  ans.reset();
                                                                   if (mask[i]) r.push_back(i);
                                                                  for (int i = 0; i < n; i++)</pre>
                                                                   d[i] = int((a[i] & mask).count());
 void add_edges( int u, bits S ) { G[ u ] = S; }
 void add_edge( int u, int v ) {
                                                                  sort(r.begin(), r.end(),
 G[u][v] = G[v][u] = 1;
                                                                   [&](int i, int j) { return d[i] > d[j]; });
                                                                  csort(r, c);
 int solve() {
                                                                  dfs(r, c, 1, mask);
                                                                  return ans; // sol[0 ~ ans-1]
  sort_by_degree(); // or simply iota( deo... )
  for ( size_t i = 0 ; i < n ; ++ i )
   deg[ i ] = G[ i ].count();
                                                               } graph;
  bits pob, nob = 0; pob.set();
                                                                3.12 Minimum Mean Cycle
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
  for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
                                                                /* minimum mean cycle O(VE) */
   size_t v = deo[ i ];
                                                                struct MMC{
   bits tmp; tmp[v] = 1;
                                                                #define FZ(n) memset((n),0,sizeof(n))
   BK( tmp, pob & G[ v ], nob & G[ v ] );
                                                                #define E 101010
   pob[ v ] = 0, nob[ v ] = 1;
                                                                #define V 1021
                                                                #define inf 1e9
                                                                 struct Edge { int v,u; double c; };
  return static_cast< int >( ans.count() );
                                                                 int n, m, prv[V][V], prve[V][V], vst[V];
};
                                                                 Edge e[E];
                                                                 vector<int> edgeID, cycle, rho;
3.11 MaxCliqueDyn
                                                                 double d[V][V];
                                                                 void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
constexpr int kN = 150;
struct MaxClique { // Maximum Clique
bitset<kN> a[kN], cs[kN];
                                                                 void add_edge( int vi , int ui , double ci )
                                                                 { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
 int ans, sol[kN], q, cur[kN], d[kN], n;
 void init(int _n) {
 n = n, and q = 0;
                                                                  for(int i=0; i<n; i++) d[0][i]=0;</pre>
 for (int i = 0; i < n; i++) a[i].reset();</pre>
                                                                  for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);
  for(int j=0; j<m; j++) {</pre>
 void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
 void csort(vector<int> &r, vector<int> &c)
                                                                    int v = e[j].v, u = e[j].u;
                                                                    if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
  d[i+1][u] = d[i][v]+e[j].c;
  int mx = 1, km = max(ans - q + 1, 1), t = 0,
m = int(r.size());
  cs[1].reset(); cs[2].reset();
                                                                     prv[i+1][u] = v;
  for (int i = 0; i < m; i++) {
  int p = r[i], k = 1;</pre>
                                                                     prve[i+1][u] = j;
   while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
                                                                  }
   cs[k][p] = 1;
   if (k < km) r[t++] = p;
                                                                 double solve(){
                                                                  // returns inf if no cycle, mmc otherwise
  c.resize(m);
                                                                  double mmc=inf;
  if(t) c[t-1] = 0;
                                                                  int st = -1;
  for (int k = km; k <= mx; k++) {</pre>
                                                                  bellman_ford();
   for (int p = int(cs[k]._Find_first());
                                                                  for(int i=0; i<n; i++) {</pre>
```

double avg=-inf;

```
for(int k=0; k<n; k++) {</pre>
                                                               int q; vector< int > G[N];
    if(d[n][i]<inf-eps)</pre>
                                                               struct Que{
     avg=max(avg,(d[n][i]-d[k][i])/(n-k));
                                                                int u, v, id;
                                                               } que[ N ];
    else avg=max(avg,inf);
                                                               int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
                                                               void dfs( int u, int f ) {
                                                                dfn[ u ] = dfn_++; int saved_rbp = stk_;
                                                                for ( int v : G[ u ] ) {
  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
                                                                 if ( v == f ) continue;
  for (int i=n; !vst[st]; st=prv[i--][st]) {
   vst[st]++;
                                                                 dfs( v, u );
                                                                 if ( stk_ - saved_rbp < SQRT_N ) continue;</pre>
   edgeID.PB(prve[i][st]);
                                                                 for ( ++ block_ ; stk_ != saved_rbp ; )
  block_id[ stk[ -- stk_ ] ] = block_;
   rho.PB(st);
  while (vst[st] != 2) {
  int v = rho.back(); rho.pop_back();
                                                                stk[ stk_ ++ ] = u;
   cycle.PB(v);
                                                               bool inPath[ N ];
   vst[v]++;
                                                               void Diff( int u ) {
                                                                if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
  reverse(ALL(edgeID));
  edgeID.resize(SZ(cycle));
                                                                else { /*add this edge*/ }
  return mmc:
                                                               void traverse( int& origin_u, int u ) {
                                                                for ( int g = lca( origin_u, u )
} mmc;
                                                                 origin_u != g ; origin_u = parent_of[ origin_u ] )
                                                                  Diff( origin_u );
3.13 Minimum Steiner Tree
                                                                for (int v = u; v != origin_u; v = parent_of[v])
// Minimum Steiner Tree
                                                                 Diff( v );
// 0(V 3^T + V^2 2^T)
                                                                origin_u = u;
struct SteinerTree {
#define V 33
                                                               void solve() {
#define T 8
                                                                dfs(1, 1);
#define INF 1023456789
                                                                while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
 int n, dst[V][V], dp[1 << T][V], tdst[V];</pre>
                                                                sort( que, que + q, [](const Que& x, const Que& y) {
 void init(int _n) {
                                                                 return tie( block_id[ x.u ], dfn[ x.v ] )
  n = _n;
                                                                     < tie( block_id[ y.u ], dfn[ y.v ] );
  for (int i = 0; i < n; i++) {</pre>
                                                                } );
   for (int j = 0; j < n; j++)
dst[i][j] = INF * (i != j);</pre>
                                                                int U = 1, V = 1;
                                                                for ( int i = 0 ; i < q ; ++ i ) {
  pass( U, que[ i ].u );
  pass( V, que[ i ].v );</pre>
 void add_edge(int ui, int vi, int wi) {
                                                                 // we could get our answer of que[ i ].id
  dst[ui][vi] = min(dst[ui][vi], wi);
  dst[vi][ui] = min(dst[vi][ui], wi);
                                                               }
                                                               /*
 void shortest_path() {
                                                               Method 2:
  for (int k = 0; k < n; k++)
                                                               dfs u:
   for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)</pre>
                                                                push u
                                                                iterate subtree
     dst[i][j] = min(dst[i][j], dst[i][k] + dst[k][j]);
                                                                push u
                                                               Let P = LCA(u, v), and St(u) \le St(v)
 int solve(const vector<int> &ter) {
                                                               if (P == u) query[St(u), St(v)]
  int t = (int)ter.size();
                                                               else query[Ed(u), St(v)], query[St(P), St(P)]
  for (int i = 1; i < (1 << t); i++)
   fill_n(dp[i], n, INF);
  fill_n(dp[0], n, 0);
                                                               3.15 Virtural Tree
  for (int msk = 1; msk < (1 << t); msk++) {</pre>
   if (msk == (msk & (-msk))) {
                                                               vector<pair<int, int>> build(vector<int> vs) { // tree
    int who = _{-}lg(msk);
    for (int i = 0; i < n; i++)
                                                                vector<pair<int, int>> res;
     dp[msk][i] = dst[ter[who]][i];
                                                                sort(vs.begin(), vs.end(), [](int i, int j) { return
                                                                dfn[i] < dfn[j]; });
vector<int> s = {0};
    continue:
   for (int i = 0; i < n; i++)</pre>
                                                                for (int v : vs) if (v != 0) {
    for (int submsk = (msk - 1) & msk; submsk; submsk =
                                                                 int o = lca(v, s.back());
     (submsk - 1) & msk)
                                                                 if (o != s.back()) {
     dp[msk][i] = min(dp[msk][i], dp[submsk][i] + dp[
                                                                  while (s.size() >= 2 and dfn[s[s.size() - 2]] >= dfn
    msk ^ submsk][i]);
                                                                   [o]) {
   for (int i = 0; i < n; i++) {</pre>
                                                                   res.emplace_back(s[s.size() - 2], s.back());
    tdst[i] = INF;
                                                                   s.pop_back();
    for (int j = 0; j < n; j++)
     tdst[i] = min(tdst[i], dp[msk][j] + dst[j][i]);
                                                                  if (s.back() != o) {
                                                                   res.emplace_back(s.back(), o);
   copy_n(tdst, n, dp[msk]);
                                                                   s.back() = o;
  }
  int ans = INF;
  for (int i = 0; i < n; i++)
                                                                 s.push_back(v);
   ans = min(ans, dp[(1 << t) - 1][i]);
  return ans;
                                                                for (size_t i = 1; i < s.size(); ++i)</pre>
                                                                 res.emplace_back(s[i - 1], s[i]);
} solver;
                                                                return res;
```

# 4 Matching & Flow

```
4.1 Bipartite Matchina
struct BipartiteMatching {
vector<int> X[N];
int fX[N], fY[N], n;
bitset<N> vis;
bool dfs(int x)
 for (auto i:X[x]) {
  if (vis[i]) continue;
  vis[i] = true;
  if (fY[i]==-1 || dfs(fY[i])){
   fY[fX[x] = i] = x;
    return true;
 return false;
void init(int n_, int m) {
 vis.reset();
 fill(X, X + (n = n_), vector<int>());
 memset(fX, -1, sizeof(int) * n);
 memset(fY, -1, sizeof(int) * m);
void add_edge(int x, int y){
 X[x].push_back(y); }
int solve() { // return how many pair matched
 int cnt = 0;
 for(int i=0;i<n;i++) {</pre>
  vis.reset();
  cnt += dfs(i);
 return cnt:
4.2 Dijkstra Cost Flow
// kN = #(vertices)
// MCMF.{Init, AddEdge, MincostMaxflow}
// MincostMaxflow(source, sink, flow_limit, &cost)
// => flow
using Pii = pair<int, int>;
constexpr int kInf = 0x3f3f3f3f, kN = 500;
struct Edge {
int to, rev, cost, flow;
struct MCMF { // 0-based
int n{}, m{}, s{}, t{};
vector<Edge> graph[kN];
// Larger range for relabeling
int64_t dis[kN] = {}, h[kN] = {};
int p[kN] = {};
void Init(int nn) {
 n = nn;
 for (int i = 0; i < n; i++) graph[i].clear();</pre>
void AddEdge(int u, int v, int f, int c) {
 graph[u].push_back({v,
  static_cast<int>(graph[v].size()), c, f});
 graph[v].push_back(
   {u, static_cast<int>(graph[u].size()) - 1,
    -c, 0});
bool Dijkstra(int &max_flow, int64_t &cost) {
 priority_queue<Pii, vector<Pii>, greater<>> pq;
 fill_n(dis, n, kInf);
 dis[s] = 0;
 pq.emplace(0, s);
 while (!pq.empty()) {
  auto u = pq.top();
  pq.pop();
   int v = u.second;
  if (dis[v] < u.first) continue;</pre>
  for (auto &e : graph[v]) {
    auto new_dis =
     dis[v] + e.cost + h[v] - h[e.to];
    if (e.flow > 0 && dis[e.to] > new_dis) {
     dis[e.to] = new_dis;
     p[e.to] = e.rev;
```

pq.emplace(dis[e.to], e.to);

```
if (dis[t] == kInf) return false;
  for (int i = 0; i < n; i++) h[i] += dis[i];</pre>
  int d = max_flow;
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   d = min(d, graph[e.to][e.rev].flow);
  max_flow -= d;
  cost += int64_t(d) * h[t];
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   e.flow += d;
   graph[e.to][e.rev].flow -= d;
  return true;
 int MincostMaxflow(
  int ss, int tt, int max_flow, int64_t &cost) {
  this->s = ss, this->t = tt;
  cost = 0;
  fill_n(h, n, 0);
  auto orig_max_flow = max_flow;
  while (Dijkstra(max_flow, cost) && max_flow) {}
  return orig_max_flow - max_flow;
};
4.3 Dinic
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
    int to, rev;
    Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
      int u = bfs.front(); bfs.pop();
      for (auto e: G[u]) {
        if (e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
        bfs.push(e.to); lv[e.to] = lv[u] + 1;
      }
    return lv[ed] != -1;
  Cap DFS(int u, Cap f){
  if (u == ed) return f;
    Cap ret = 0;
    for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
      auto &e = G[u][i];
      if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
      Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
      G[e.to][e.rev].cap += nf;
      if (f == 0) return ret;
    if (ret == 0) lv[u] = -1;
    return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
    G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
    st = st_, ed = ed_; Cap ret = 0;
    while (BFS()) {
      idx.assign(n, 0);
      Cap f = DFS(st, numeric_limits<Cap>::max());
      ret += f;
      if (f == 0) break;
```

```
return ret;
};
```

#### 4.4 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source  ${\cal S}$  and sink  ${\cal T}$ .
  - 2. For each edge (x,y,l,u), connect  $x \to y$  with capacity u-l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If  $in(v)>ar{0}$  , connect S o v with capacity in(v) , otherwise, connect  $v \to T$  with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f 
      eq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the
    - maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect t o s with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f'\neq \sum_{v\in V, in(v)>0}in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e \, + \, f_e$  , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph(X, Y)
  - 1. Redirect every edge:  $y \to x$  if  $(x, y) \in M$ ,  $x \to y$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- Minimum cost cyclic flow
  - 1. Consruct super source  ${\cal S}$  and sink  ${\cal T}$
  - 2. For each edge (x,y,c), connect x o y with (cost,cap) = (c,1) if c>0, otherwise connect  $y\to x$  with (cost, cap)=(-c,1)
  - 3. For each edge with  $c\,<\,0$ , sum these cost as K, then increase d(y)by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v)>0, connect S o v with (cost, cap)=(0, d(v))
  - 5. For each vertex v with d(v) < 0, connect  $v \to T$  with (cost, cap) =(0, -d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let K be the sum of all weights
  - 3. Connect source s o v,  $v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity
  - 5. For  $v \in {\it G}$ , connect it with sink  $v \to t$  with capacity K + 2T - $(\sum_{e \in E(v)} w(e)) - 2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight w(u,v).
  - $ightarrow \ v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
  - 1. If  $p_v>0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge
  - (v,t) with capacity  $-p_v$ . 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
  - 3. The mincut is equivalent to the maximum profit of a subset of projects
- 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity  $c_x$  and create edge (s,y) with capacity  $c_i$
- 2. Create edge (x,y) with capacity  $c_{xy}$ .

  3. Create edge (x,y) and edge (x',y') with capacity  $c_{xyx'y'}$ .

# General Graph Matching

```
namespace matching {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
g[u].push_back(v);
g[v].push_back(u);
int Find(int u) {
```

```
return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
 static int tk = 0; tk++;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
void Blossom(int x, int y, int 1) {
 while (Find(x) != 1) {
  pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
if (fa[x] == x) fa[x] = 1;
  if (fa[y] == y) fa[y] = 1;
  x = pre[y];
bool Bfs(int r, int n) {
 for (int i = 0; i \le n; ++i) fa[i] = i, s[i] = -1;
 while (!q.empty()) q.pop();
 q.push(r);
 s[r] = 0:
 while (!q.empty()) {
  int x = q.front(); q.pop();
for (int u : g[x]) {
   if (s[u] == -1) {
    pre[u] = x, s[u] = 1;
    if (match[u] == n) {
     for (int a = u, b = x, last; b != n; a = last, b =
     pre[a])
      last = match[b], match[b] = a, match[a] = b;
     return true;
    q.push(match[u]);
    s[match[u]] = 0;
   } else if (!s[u] && Find(u) != Find(x)) {
    int 1 = LCA(u, x, n);
Blossom(x, u, 1);
    Blossom(u, x, 1);
  }
 return false;
int Solve(int n) {
 int res = 0;
 for (int x = 0; x < n; ++x) {
  if (match[x] == n) res += Bfs(x, n);
 return res;
}}
4.6 Global Min-Cut
```

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
pair<int, int> phase(int n) {
memset(v, false, sizeof(v));
 memset(g, 0, sizeof(g));
 int s = -1, t = -1;
 while (true) {
  int c = -1;
  for (int i = 0; i < n; ++i) {
   if (del[i] || v[i]) continue;
   if (c == -1 \mid | g[i] > g[c]) c = i;
  if (c == -1) break;
  v[s = t, t = c] = true;
  for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
   g[i] += w[c][i];
 return make_pair(s, t);
```

w.assign(n, vector<lld>(n));

```
slk.resize(n); pre.resize(n);
int mincut(int n) {
                                                                vl.resize(n); vr.resize(n);
int cut = 1e9;
                                                               void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 memset(del, false, sizeof(del));
 for (int i = 0; i < n - 1; ++i) {
                                                               11d solve() {
                                                                for (int i = 0; i < n; ++i)
  int s, t; tie(s, t) = phase(n);
  del[t] = true; cut = min(cut, g[t]);
for (int j = 0; j < n; ++j) {</pre>
                                                                 hl[i] = *max_element(w[i].begin(), w[i].end());
                                                                for (int i = 0; i < n; ++i) bfs(i);</pre>
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
                                                                11d res = 0;
  }
                                                                for (int i = 0; i < n; ++i) res += w[i][f1[i]];</pre>
                                                                return res;
return cut;
                                                              } km;
4.7 GomoryHu Tree
                                                              4.9 Minimum Cost Circulation
int g[maxn];
                                                              struct Edge { int to, cap, rev, cost; };
                                                              vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
vector<edge> GomoryHu(int n){
 vector<edge> rt;
 for(int i=1;i<=n;++i)g[i]=1;</pre>
                                                              bool mark[kN];
 for(int i=2;i<=n;++i){</pre>
                                                              int NegativeCycle(int n) {
                                                               memset(mark, false, sizeof(mark));
  int t=g[i];
  flow.reset(); // clear flows on all edge
                                                               memset(dist, 0, sizeof(dist));
 rt.push_back({i,t,flow(i,t)});
flow.walk(i); // bfs points that connected to i (use
                                                               int upd = -1;
                                                               for (int i = 0; i <= n; ++i) {
    edges not fully flow)
                                                                for (int j = 0; j < n; ++j) {
                                                                 int idx = 0;
  for(int j=i+1;j<=n;++j){</pre>
                                                                 for (auto &e : g[j])
  if(g[j]==t && flow.connect(j))g[j]=i; // check if i
    can reach j
                                                                  if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
  }
                                                                    dist[e.to] = dist[j] + e.cost;
                                                                    pv[e.to] = j, ed[e.to] = idx;
                                                                    if (i == n) {
 return rt;
                                                                    upd = j;
                                                                     while(!mark[upd])mark[upd]=1,upd=pv[upd];
4.8 Kuhn Munkres
                                                                     return upd;
class KM {
                                                                    }
private:
 static constexpr lld INF = 1LL << 60;</pre>
                                                                   idx++;
 vector<lld> h1,hr,slk;
 vector<int> f1,fr,pre,qu;
 vector<vector<lld>> w;
 vector<bool> v1,vr;
                                                               return -1;
 int n, ql, qr;
 bool check(int x) {
                                                              int Solve(int n) {
  if (v1[x] = true, f1[x] != -1)
                                                               int rt = -1, ans = 0;
   return vr[qu[qr++] = f1[x]] = true;
                                                               while ((rt = NegativeCycle(n)) >= 0) {
  while (x != -1) swap(x, fr[fl[x] = pre[x]]);
                                                                memset(mark, false, sizeof(mark));
                                                                vector<pair<int, int>> cyc;
  return false:
                                                                while (!mark[rt]) {
 void bfs(int s) {
                                                                 cyc.emplace_back(pv[rt], ed[rt]);
  fill(slk.begin(), slk.end(), INF);
                                                                 mark[rt] = true;
  fill(v1.begin(), v1.end(), false);
fill(vr.begin(), vr.end(), false);
                                                                 rt = pv[rt];
                                                                reverse(cyc.begin(), cyc.end());
  ql = qr = 0;
  vr[qu[qr++] = s] = true;
                                                                int cap = kInf;
  while (true) {
                                                                for (auto &i : cyc)
                                                                 auto &e = g[i.first][i.second];
   11d d:
   while (ql < qr) {</pre>
                                                                 cap = min(cap, e.cap);
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
     if(!v1[x]&&slk[x]>=(d=h1[x]+hr[y]-w[x][y])){
                                                                for (auto &i : cyc)
      if (pre[x] = y, d) slk[x] = d;
                                                                 auto &e = g[i.first][i.second];
                                                                 e.cap -= cap;
      else if (!check(x)) return;
                                                                 g[e.to][e.rev].cap += cap;
     }
                                                                 ans += e.cost * cap;
                                                                }
   d = INF;
   for (int x = 0; x < n; ++x)
                                                               return ans;
    if (!v1[x] \&\& d > s1k[x]) d = s1k[x];
   for (int x = 0; x < n; ++x) {
                                                                     Minimum Cost Maximum Flow
                                                              4.10
    if (vl[x]) hl[x] += d;
                                                              class MiniCostMaxiFlow{
    else slk[x] -= d;
    if (vr[x]) hr[x] -= d;
                                                               using Cap = int; using Wei = int64_t;
                                                               using PCW = pair<Cap,Wei>;
   for (int x = 0; x < n; ++x)
                                                               static constexpr Cap INF_CAP = 1 << 30;</pre>
    if (!v1[x] && !slk[x] && !check(x)) return;
                                                               static constexpr Wei INF_WEI = 1LL<<60;</pre>
  }
                                                              private:
                                                               struct Edge{
public:
                                                                int to, back;
                                                                Cap cap; Wei wei;
 void init( int n_ ) {
  qu.resize(n = n_);
                                                                Edge() {}
 fl.assign(n, -1); fr.assign(n, -1);
hr.assign(n, 0); hl.resize(n);
                                                                Edge(int a,int b, Cap c, Wei d):
                                                                 to(a),back(b),cap(c),wei(d) {}
```

};

] = u; }

```
int ori, edd;
                                                             void set_slack(int x) {
vector<vector<Edge>> G;
                                                              slack[x] = 0;
vector<int> fa, wh;
                                                              for (int u = 1; u <= n; ++u)</pre>
vector<bool> inq;
                                                               if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
vector<Wei> dis;
                                                                update_slack(u, x);
PCW SPFA(){
 fill(inq.begin(),inq.end(),false);
                                                             void q_push(int x)
 fill(dis.begin(), dis.end(), INF_WEI);
                                                              if (x \le n) q.push(x);
  queue<int> qq; qq.push(ori);
                                                              else for (size_t i = 0; i < flo[x].size(); i++)</pre>
 dis[ori] = 0;
                                                                 q_push(flo[x][i]);
 while(not qq.empty()){
  int u=qq.front();qq.pop();
                                                             void set_st(int x, int b) {
   inq[u] = false;
                                                              st[x] = b;
   for(int i=0;i<SZ(G[u]);++i){</pre>
                                                              if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)</pre>
   Edge e=G[u][i];
                                                                  set_st(flo[x][i], b);
    int v=e.to; Wei d=e.wei;
                                                             int get_pr(int b, int xr) {
    if(e.cap <= 0 | |dis[v] <= dis[u] + d)
                                                              int pr = find(flo[b].begin(), flo[b].end(), xr) - flo
    continue
    dis[v] = dis[u] + d;
                                                                 [b].begin();
    fa[v] = u, wh[v] = i;
                                                              if (pr % 2 == 1) {
   if (inq[v]) continue;
                                                               reverse(flo[b].begin() + 1, flo[b].end());
    qq.push(v);
                                                                return (int)flo[b].size() - pr;
    inq[v] = true;
                                                              return pr;
  if(dis[edd]==INF_WEI) return {-1, -1};
                                                             void set_match(int u, int v) {
                                                              match[u] = g[u][v].v;
 Cap mw=INF_CAP;
 for(int i=edd;i!=ori;i=fa[i])
                                                              if (u <= n) return;</pre>
                                                              edge e = g[u][v];
int xr = flo_from[u][e.u], pr = get_pr(u, xr)
  mw=min(mw,G[fa[i]][wh[i]].cap);
  for (int i=edd;i!=ori;i=fa[i]){
  auto &eg=G[fa[i]][wh[i]];
                                                              for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo</pre>
                                                                 [u][i ^ 1]);
   eq.cap -= mw:
  G[eg.to][eg.back].cap+=mw;
                                                              set_match(xr,
                                                              rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
 return {mw, dis[edd]};
                                                                 end());
                                                             void augment(int u, int v) {
public:
void init(int n){
                                                              for (; ; )
 G.clear();G.resize(n);
                                                               int xnv = st[match[u]];
                                                               set_match(u, v);
 fa.resize(n);wh.resize(n);
 inq.resize(n); dis.resize(n);
                                                                if (!xnv) return;
                                                               set_match(xnv, st[pa[xnv]]);
                                                               u = st[pa[xnv]], v = xnv;
void add_edge(int st, int ed, Cap c, Wei w){
 G[st].emplace_back(ed,SZ(G[ed]),c,w);
 G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
                                                             int get_lca(int u, int v) {
PCW solve(int a, int b){
                                                              static int t = 0;
 ori = a, edd = b;
                                                              for (++t; u || v; swap(u, v)) {
  if (u == 0) continue;
 Cap cc=0; Wei ww=0;
                                                               if (vis[u] == t) return u;
 while(true){
  PCW ret=SPFA();
                                                               vis[u] = t:
   if(ret.first==-1) break;
                                                               u = st[match[u]]
                                                               if (u) u = st[pa[u]];
  cc+=ret.first;
  ww+=ret.first * ret.second;
                                                              }
                                                              return 0;
 return {cc,ww};
}
                                                             void add_blossom(int u, int lca, int v) {
} mcmf;
                                                              int b = n + 1;
                                                              while (b \le n_x \& st[b]) ++b;
      Maximum Weight Graph Matching
                                                              if (b > n_x) ++n_x;
struct WeightGraph {
                                                              lab[b] = 0, S[b] = 0
                                                              match[b] = match[lca];
static const int inf = INT_MAX;
                                                              flo[b].clear();
static const int maxn = 514;
                                                              flo[b].push_back(lca);
struct edge {
                                                              for (int x = u, y; x != lca; x = st[pa[y]])
 int u, v, w;
                                                               flo[b].push_back(x), flo[b].push_back(y = st[match[x
 edge(){}
 edge(int u, int v, int w): u(u), v(v), w(w) {}
                                                                 ]]), q_push(y);
                                                               reverse(flo[b].begin() + 1, flo[b].end())
                                                               for (int x = v, y; x != lca; x = st[pa[y]])
int n, n_x;
edge g[maxn * 2][maxn * 2];
                                                               flo[b].push_back(x), flo[b].push_back(y = st[match[x
int lab[maxn * 2];
                                                                 ]]), q_push(y);
                                                              set_st(b, b);
int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa
                                                              for (int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].w
int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
                                                              for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
    maxn * 2]
vector<int> flo[maxn * 2];
                                                              for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
queue<int> q;
                                                               int xs = flo[b][i];
                                                                for (int x = 1; x <= n_x; ++x)
int e_delta(const edge &e) { return lab[e.u] + lab[e.v
    ] - g[e.u][e.v].w * 2; }
                                                                if (g[b][x].w == 0 \mid \mid e_delta(g[xs][x]) < e_delta(g
void update_slack(int u, int x) { if (!slack[x] ||
                                                                 [b][x]))
    e_delta(g[u][x]) < e_delta(g[slack[x]][x])) slack[x]
                                                                 g[b][x] = g[xs][x], g[x][b] = g[x][xs];
```

for (int x = 1; x <= n; ++x)

```
if (flo_from[xs][x]) flo_from[b][x] = xs;
 set_slack(b);
void expand_blossom(int b) {
 for (size_t i = 0; i < flo[b].size(); ++i)
set_st(flo[b][i], flo[b][i]);</pre>
 int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
 for (int i = 0; i < pr; i += 2) {
  int xs = flo[b][i], xns = flo[b][i + 1];</pre>
  pa[xs] = g[xns][xs].u;
  S[xs] = 1, S[xns] = 0;
slack[xs] = 0, set_slack(xns);
  q_push(xns);
 S[xr] = 1, pa[xr] = pa[b];
 for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  S[xs] = -1, set_slack(xs);
 st[b] = 0;
bool on_found_edge(const_edge &e) {
 int u = st[e.u], v = st[e.v];
 if (S[v] == -1) {
pa[v] = e.u, S[v] = 1;
  int nu = st[match[v]];
  slack[v] = slack[nu] = 0;
  S[nu] = 0, q_push(nu);
 } else if (S[v] == 0) {
  int lca = get_lca(u, v);
  if (!lca) return augment(u,v), augment(v,u), true;
  else add_blossom(u, lca, v);
 return false;
bool matching() {
 memset(S + 1, -1, sizeof(int) * n_x);
 memset(slack + 1, 0, sizeof(int) * n_x);
 q = queue<int>();
 for (int x = 1; x <= n_x; ++x)
  if (st[x] == x \&\& !match[x]) pa[x] = 0, S[x] = 0,
   q_push(x);
 if (q.empty()) return false;
 for (; ; )
  while (q.size()) {
   int u = q.front(); q.pop();
   if (S[st[u]] == 1) continue
   for (int v = 1; v <= n; ++v)
    if (g[u][v].w > 0 && st[u] != st[v]) {
     if (e_delta(g[u][v]) == 0) {
      if (on_found_edge(g[u][v])) return true;
     } else update_slack(u, st[v]);
  int d = inf;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b && S[b] == 1) d = min(d, lab[b] / 2)
  for (int x = 1; x <= n_x; ++x)
if (st[x] == x && slack[x]) {</pre>
    if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x])
    else if (S[x] == 0) d = min(d, e_delta(g[slack[x
   ]][x]) / 2);
  for (int u = 1; u <= n; ++u) {
   if (S[st[u]] == 0) {
    if (lab[u] <= d) return 0;</pre>
    lab[u] -= d;
   } else if (S[st[u]] == 1) lab[u] += d;
  for (int b = n + 1; b <= n_x; ++b)
   if (st[b] == b) {
    if (S[st[b]] == 0) lab[b] += d * 2;
    else if (S[st[b]] == 1) lab[b] -= d * 2;
   }
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
   if (st[x] == x && slack[x] && st[slack[x]] != x &&
   e_delta(g[slack[x]][x]) == 0)
```

```
if (on_found_edge(g[slack[x]][x])) return true;
   for (int b = n + 1; b \le n_x; ++b)
    if (st[b] == b && S[b] == 1 && lab[b] == 0)
    expand_blossom(b);
  return false;
 pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n_x = n;
  int n_matches = 0;
  long long tot_weight = 0;
  for (int u = 0; u \le n; ++u) st[u] = u, flo[u].clear
  int w_max = 0;
  for (int u = 1; u <= n; ++u)
   for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);</pre>
    w_{max} = max(w_{max}, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
  while (matching()) ++n_matches;
  for (int u = 1; u <= n; ++u)
   if (match[u] && match[u] < u)</pre>
    tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
 void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
    g[vi][ui].w = wi; }
 void init(int _n) {
  n = _n;
  for (int u = 1; u <= n; ++u)
   for (int v = 1; v <= n; ++v)
    g[u][v] = edge(u, v, 0);
};
```

### 5 Math

# 5.1 $\lfloor \frac{n}{i} \rfloor$ Enumeration

 $T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor$ 

#### 5.2 Strling Number

# 5.2.1 First Kind

 $S_1(n,k)$  counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1)\dots(x+n-1) = \sum_{k=0}^n S_1(n,k)x^k$$

$$g(x) = x(x+1)\dots(x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

#### 5.2.2 Second Kind

 $S_2(n,k)$  counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

#### 5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b-=(x/y)*a;
}
```

# 5.4 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
vector<T> d(output.size() + 1), me, he;
for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
  for (size_t j = 0; j < me.size(); ++j)
d[i] += output[i - j - 2] * me[j];
if ((d[i] -= output[i - 1]) == 0) continue;</pre>
  if (me.empty()) {
   me.resize(f = i);
   continue;
  vector<T> o(i - f - 1);
 T k = -d[i] / d[f]; o.push_back(-k);
  for (T x : he) o.push_back(x * k);
 if (o.size() < me.size()) o.resize(me.size());</pre>
  for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
  if (i-f+he.size() >= me.size()) he = me, f = i;
 me = o;
return me;
```

```
5.5 Charateristic Polynomial
vector<vector<int>> Hessenberg(const vector<vector<int
    >> &A) {
int N = A.size();
vector<vector<int>> H = A;
for (int i = 0; i < N - 2; ++i) {
 if (!H[i + 1][i]) {
  for (int j = i + 2; j < N; ++j) {
    if (H[j][i]) {
    for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
    ][k]);
     for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k
    ][j]);
    break:
 if (!H[i + 1][i]) continue;
 int val = fpow(H[i + 1][i], kP - 2);
 for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;
  for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
    * H[i + 1][k] * (kP - coef)) % kP;
   for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
}
return H;
vector<int> CharacteristicPoly(const vector<vector<int
    >> &A) {
int N = A.size();
auto H = Hessenberg(A);
for (int i = 0; i < N; ++i) {
 for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
vector<vector<int>>> P(N + 1, vector<int>(N + 1));
P[0][0] = 1;
for (int i = 1; i <= N; ++i) {
 P[i][0] = 0;
 for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
 int val = 1:
  for (int j = i - 1; j >= 0; --j)
  int coef = 1LL * val * H[j][i
                                  - 1] % kP;
  for (int k = 0; k <= j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
  if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 }
if (N & 1) {
 for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
return P[N];
```

#### 5.6 Chinese Remainder

```
|lld crt(lld ans[], lld pri[], int n){
```

```
11d M = 1, ret = 0;
 for(int i=0;i<n;i++) M *= pri[i];</pre>
 for(int i=0;i<n;i++){</pre>
  lld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
  ret += (ans[i]*(M/pri[i])%M * iv)%M;
  ret %= M;
 return ret;
/*
Another:
x = a1 \% m1
x = a2 \% m2
g = gcd(m1, m2)
assert((a1-a2)\%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
0 <= x < lcm(m1, m2)
*/
5.7 De-Bruijn
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
  if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
 } else {
  aux[t] = aux[t - p];
  db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {
   aux[t] = i;
   db(t + 1, t, n, k);
int de_bruijn(int k, int n) {
 // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
 if (k == 1) {
  res[0] = 0:
  return 1;
 for (int i = 0; i < k * n; i++) aux[i] = 0;
 sz = 0;
 db(1, 1, n, k);
 return sz;
5.8
      DiscreteLog
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
  // x^? \neq M \pmod{M}
  Int t = 1, c = 0, g = 1;
  for (Int M_ = M; M_ > 0; M_ >>= 1)
```

```
template < typename Int>
Int BSGS(Int x, Int y, Int M) {
    // x^? \equiv y (mod M)
    Int t = 1, c = 0, g = 1;
    for (Int M_ = M; M_ > 0; M_ >>= 1)
        g = g * x % M;
    for (g = gcd(g, M); t % g != 0; ++c) {
        if (t == y) return c;
        t = t * x % M;
    }
    if (y % g != 0) return -1;
    t /= g, y /= g, M /= g;
    Int h = 0, gs = 1;
    for (; h * h < M; ++h) gs = gs * x % M;
    unordered_map<Int, Int> bs;
    for (Int s = 0; s < h; bs[y] = ++s)
        y = y * x % M;
    for (Int s = 0; s < M; s += h) {
        t = t * gs % M;
        if (bs.count(t)) return c + s + h - bs[t];
    }
    return -1;
}</pre>
```

#### 5.9 Extended Euler

```
a^b \equiv \begin{cases} a^b \mod \varphi(m) + \varphi(m) & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m
```

#### 5.10 ExtendedFloorSum

```
g(a, b, c, n) = \sum_{i=0}^{n} i \lfloor \frac{ai + b}{c} \rfloor
                            \left( \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \right)
                                                                                                           a \geq c \vee b \geq c
                                                                                                           n<0\vee a=0
                             \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
                            -h(c, c-b-1, a, m-1)),
                                                                                                           otherwise
h(a,b,c,n) = \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2
                           \left( \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right)
                             + \left| \frac{a}{c} \right| \cdot \left| \frac{b}{c} \right| \cdot n(n+1)
                             +h(a \bmod c, b \bmod c, c, n)
                            +2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)
                             +2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
                                                                                                           a \ge c \lor b \ge c
                                                                                                           n<0\vee a=0
                             nm(m+1) - 2g(c, c-b-1, a, m-1)
                            -2f(c, c-b-1, a, m-1) - f(a, b, c, n), otherwise
```

```
Fast Fourier Transform
const int mod = 1000000007:
const int M1 = 985661441; // G = 3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
  static_assert (M1 <= M2 && M2 <= M3);
  constexpr int64_t r12 = modpow(M1, M2-2, M2);
 constexpr int64_t r13 = modpow(M1, M3-2, M3);
 constexpr int64_t r23 = modpow(M2, M3-2, M3);
  constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
 B = (B - A + M2) * r12 % M2;
 C = (C - A + M3) * r13 % M3;
 C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
 for (int i = 0; i <= maxn; i++)</pre>
  omega[i] = cplx(cos(2 * pi * j / maxn),
     sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
 int z = __builtin_ctz(n) - 1;
 for (int i = 0; i < n; ++i) {
 int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^{=(i >> j & 1) << (z - j);
  if (x > i) swap(v[x], v[i]);
 for (int s = 2; s <= n; s <<= 1) {
  int z = s >> 1;
  for (int i = 0; i < n; i += s) {
  for (int k = 0; k < z; ++k) {
  cplx x = v[i + z + k] * omega[maxn / s * k];
  v[i + z + k] = v[i + k] - x;</pre>
    v[i+k] = v[i+k] + x;
   }
void ifft(vector<cplx> &v, int n) {
fft(v, n); reverse(v.begin() + 1, v.end());
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
VL convolution(const VI &a, const VI &b) {
 // Should be able to handle N <= 10^5, C <= 10^4
 int sz = 1;
 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 vector<cplx> v(sz);
 for (int i = 0; i < sz; ++i) {
  double re = i < a.size() ? a[i] : 0;
  double im = i < b.size() ? b[i] : 0;</pre>
  v[i] = cplx(re, im);
fft(v, sz);
```

```
for (int i = 0; i <= sz / 2; ++i) {
  int j = (sz - i) & (sz - 1);
  cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
  * cplx(0, -0.25);
if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
     ].conj()) * cplx(0, -0.25);
  v[i] = x;
 ifft(v, sz);
 VL c(sz);
 for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
 return c:
VI convolution_mod(const VI &a, const VI &b, int p) {
 int sz = 1;
 while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
 vector<cplx> fa(sz), fb(sz);
 for (int i = 0; i < (int)a.size(); ++i)</pre>
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 for (int i = 0; i < (int)b.size(); ++i)</pre>
  fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
 fft(fa, sz), fft(fb, sz);
 double r = 0.25 / sz;
 cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
  int j = (sz - i) & (sz - 1);
  cplx a1 = (fa[i] + fa[j].conj());
  cplx a2 = (fa[i] - fa[j].conj()) * r2;
  cplx b1 = (fb[i] + fb[j].conj()) * r3;
cplx b2 = (fb[i] - fb[j].conj()) * r4;
  if (i != j) {
   cplx c1 = (fa[j] + fa[i].conj());
cplx c2 = (fa[j] - fa[i].conj()) * r2;
   cplx d1 = (fb[j] + fb[i].conj()) * r3;
   cplx d2 = (fb[j] - fb[i].conj()) * r4;
   fa[i] = c1 * d1 + c2 * d2 * r5;
   fb[i] = c1 * d2 + c2 * d1;
  fa[j] = a1 * b1 + a2 * b2 * r5;
  fb[j] = a1 * b2 + a2 * b1;
 fft(fa, sz), fft(fb, sz);
 vector<int> res(sz);
 for (int i = 0; i < sz; ++i) {
  long long a = round(fa[i].re), b = round(fb[i].re),
        c = round(fa[i].im)
  res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
 }
 return res;
}}
5.12 FloorSum
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
// @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
 llu ans = 0;
 while (true)
  if (a >= m) {
   ans += n * (n - 1) / 2 * (a / m); a %= m;
  if (b >= m) {
   ans += n * (b / m); b %= m;
  llu y_max = a * n + b;
  if (y_max < m) break;</pre>
  // y_max < m * (n + 1)
// floor(y_max / m) <= n
  n = (11u)(y_max / m), b = (11u)(y_max % m);
  swap(m, a);
 return ans;
11d floor_sum(lld n, lld m, lld a, lld b) {
 11u ans = 0;
 if (a < 0) {
  llu a2 = (a % m + m) % m;
ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
  a = a2;
 if (b < 0) {
```

11u b2 = (b % m + m) % m;

```
ans -= 1ULL * n * ((b2 - b) / m);
                                                                 for (int i = 0; i < n; i += s * 2) {
                                                                  for (int j = 0; j < s; j++) {
 b = b2;
                                                                   int a = F[i+j]
                                                                   int b = modmul(F[i+j+s], roots[s+j]);
return ans + floor_sum_unsigned(n, m, a, b);
                                                                   F[i+j] = modadd(a, b); // a + b
                                                                   F[i+j+s] = modsub(a, b); // a - b
5.13 FWT
                                                                 }
/* xor convolution:
* x = (x0, x1) , y = (y0, y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                if (inv) {
                                                                 int invn = modinv(n);
                                                                 for (int i = 0; i < n; i++)
* x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
                                                                 F[i] = modmul(F[i], invn);
                                                                 reverse(F + 1, F + n);
*z = (1/2) *z'
 * or convolution:
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
* and convolution:
                                                             NTT<2013265921, 31, 1048576> ntt;
* x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */const LL MOD = 1e9+7;
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
                                                                    Partition Number
for( int d = 1 ; d < N ; d <<= 1 ) {
                                                             int b = sqrt(n);
  int d2 = d<<1;
                                                              ans[0] = tmp[0] = 1;
  for( int s = 0 ; s < N ; s += d2 )
                                                              for (int i = 1; i <= b; i++) {
   for (int rep = 0; rep < 2; rep++)</pre>
                                                                for (int j = i; j <= n - i * i; j++)
   x[ i ] = ta+tb;
                                                                modadd(tmp[j], tmp[j-i]);
   x[ j ] = ta-tb;
                                                               for (int j = i * i; j <= n; j++)
    if( x[ i ] >= MOD ) x[ i ] -= MOD;
                                                                modadd(ans[j], tmp[j - i * i]);
    if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
                                                                    Pi Count (Linear Sieve)
if( inv )
                                                              5.17
 for( int i = 0 ; i < N ; i++ ) {</pre>
                                                              static_constexpr int N = 1000000 + 5;
  x[ i ] *= inv( N, MOD );
x[ i ] %= MOD;
                                                             1ld pi[N];
                                                             vector<int> primes;
                                                             bool sieved[N]
                                                              11d cube_root(11d x){
                                                              lld s=cbrt(x-static_cast<long double>(0.1));
      Miller Rabin
5.14
                                                               while(s*s*s <= x) ++s;
bool isprime(llu x) -
                                                               return s-1;
static auto witn = [](llu a, llu u, llu n, int t) {
                                                              11d square_root(11d x){
 if (!a) return false;
  while (t--) {
                                                              lld s=sqrt(x-static_cast<long double>(0.1));
  11u a2 = mmul(a, a, n);
                                                               while(s*s <= x) ++s;
  if (a2 == 1 && a != 1 && a != n - 1) return true;
                                                               return s-1;
  a = a2;
 }
                                                              void init(){
 return a != 1;
                                                              primes.reserve(N);
};
                                                               primes.push_back(1);
if (x < 2) return false;</pre>
                                                               for(int i=2;i<N;i++) {</pre>
if (!(x & 1)) return x == 2;
                                                                if(!sieved[i]) primes.push_back(i);
int t = __builtin_ctzll(x - 1);
llu odd = (x - 1) >> t;
                                                                pi[i] = !sieved[i] + pi[i-1];
                                                                for(int p: primes) if(p > 1) {
  if(p * i >= N) break;
for (llu m:
  {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
                                                                 sieved[p * i] = true;
  if (witn(mpow(m % x, odd, x), odd, x, t))
                                                                 if(p % i == 0) break;
   return false;
return true;
                                                              11d phi(11d m, 11d n) {
5.15 NTT
                                                               static constexpr int MM = 80000, NN = 500;
                                                               static lld val[MM][NN];
template <int mod, int G, int maxn>
                                                               if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
struct NTT {
                                                               if(n == 0) return m;
static_assert (maxn == (maxn & -maxn));
                                                               if(primes[n] >= m) return 1;
 int roots[maxn];
                                                               1ld ret = phi(m,n-1)-phi(m/primes[n],n-1);
NTT () {
                                                               if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
  int r = modpow(G, (mod - 1) / maxn);
 for (int i = maxn >> 1; i; i >>= 1) {
                                                               return ret;
  roots[i] = 1;
                                                             1ld pi_count(1ld);
  for (int j = 1; j < i;
                                                             11d P2(11d m, 11d n) {
   roots[i + j] = modmul(roots[i + j - 1], r);
                                                               lld sm = square_root(m), ret = 0;
for(lld i = n+1;primes[i]<=sm;i++)</pre>
   r = modmul(r, r);
 }
                                                                ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
                                                               return ret;
 // n must be 2^k, and 0 \le F[i] < mod
void operator()(int F[], int n, bool inv = false) {
                                                              11d pi_count(11d m) {
 for (int i = 0, j = 0; i < n; i++) {
                                                               if(m < N) return pi[m];</pre>
  if (i < j) swap(F[i], F[j]);</pre>
                                                               11d n = pi_count(cube_root(m));
   for (int k = n > 1; (j^k < k; k > = 1);
                                                               return phi(m, n) + n - 1 - P2(m, n);
  for (int s = 1; s < n; s *= 2) {
```

// does not work when n is prime

return any non-trivial factor

#### 5.18 Pollard Rho

```
llu pollard_rho(llu n) {
 static auto f = [](llu x, llu k, llu m) {
 return add(k, mul(x, x, m), m); };
if (!(n & 1)) return 2;
 mt19937 rnd(120821011);
 while (true) {
   llu y = 2, yy = y, x = rnd() % n, t = 1;
   for (llu sz = 2; t == 1; sz <<= 1, y = yy) {
     for (llu i = 0; t == 1 && i < sz; ++i) {
       yy = f(yy, x, n);
       t = gcd(yy > y ? yy - y : y - yy, n);
     }
   if (t != 1 && t != n) return t;
           Polynomial Operations
using V = vector<int>
#define fi(1, r) for (int i = int(1); i < int(r); ++i)
template <int mod, int G, int maxn> struct Poly : V {
 static uint32_t n2k(uint32_t n) {
   if (n <= 1) return 1;
   return 1u << (32 - __builtin_clz(n - 1));</pre>
 static NTT<mod,G,maxn> ntt; // coefficients in [0, P)
 explicit Poly(int n = 1) : V(n) {}
 Poly(const V &v) : V(v) {}
 Poly(const Poly &p, size_t n) : V(n) {
   copy_n(p.data(), min(p.size(), n), data());
 Poly &irev() { return reverse(data(), data() + size())
          *this; }
 Poly &isz(int sz) { return resize(sz), *this; }
 Poly &iadd(const Poly &rhs) { // n() == rhs.n()
   fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
   return *this;
 Poly &imul(int k) {
   fi(0, size())(*this)[i] = modmul((*this)[i], k);
   return *this;
 Poly Mul(const Poly &rhs) const {
   const int sz = n2k(size() + rhs.size() - 1);
   Poly X(*this, sz), Y(rhs, sz);
   ntt(X.data(), sz), ntt(Y.data(), sz);
   fi(0, sz) X[i] = modmul(X[i], Y[i]);
   ntt(X.data(), sz, true);
   return X.isz(size() + rhs.size() - 1);
 Poly Inv() const { // coef[0] != 0
   if (size() == 1) return V{modinv(*begin())};
   const int sz = n2k(size() * 2);
   Poly X = Poly(*this, (size() + 1) / 2).Inv().isz(sz),
         Y(*this, sz);
   ntt(X.data(), sz), ntt(Y.data(), sz);
   fi(0, sz) X[i] = modmul(X[i], modsub(2, modsub(2
       Y[i])));
   ntt(X.data(), sz, true);
   return X.isz(size());
 Poly Sqrt() const { // coef[0] \in [1, mod)^2
   if (size() == 1) return V{QuadraticResidue((*this)
       [0], mod)};
   Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
       size());
   return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
       + 1);
 pair<Poly, Poly> DivMod(const Poly &rhs) const {
   if (size() < rhs.size()) return {V{0}, *this};</pre>
   const int sz = size() - rhs.size() + 1;
   Poly X(rhs); X.irev().isz(sz);
Poly Y(*this); Y.irev().isz(sz);
   Poly Q = Y.Mul(X.Inv()).isz(sz).irev();
   X = rhs.Mul(Q), Y = *this;
fi(0, size()) Y[i] = modsub(Y[i], X[i]);
   return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
```

```
Poly Dx() const {
  Poly ret(size() - 1);
  fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
    1]);
  return ret.isz(max<int>(1, ret.size()));
 Poly Sx() const {
  Poly ret(size() + 1);
  fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
    this)[i]);
  return ret;
 Poly Ln() const \{ // coef[0] == 1 \}
 return Dx().Mul(Inv()).Sx().isz(size());
 Poly Exp() const \{ // coef[0] == 0 \}
  if (size() == 1) return V{1};
  Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size)
    ());
  Poly Y = X.Ln(); Y[0] = mod - 1;
  fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
  return X.Mul(Y).isz(size());
 Poly Pow(const string &K) const {
  int nz = 0;
  while (nz < size() && !(*this)[nz]) ++nz;</pre>
  int nk = 0, nk2 = 0;
  for (char c : K) {
   nk = (nk * 10 + c - '0') % mod;
   nk2 = nk2 * 10 + c - '0';
   if (nk2 * nz >= size())
    return Poly(size());
   nk2 %= mod - 1;
  if (!nk && !nk2) return Poly(V{1}, size());
  Poly X = V(data() + nz, data() + size() - nz * (nk2 -
     1));
  int x0 = X[0]
  return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
    modpow(x0, nk2)).irev().isz(size()).irev();
 Poly InvMod(int L) { // (to evaluate linear recursion)
  Poly R{1, 0}; // *this * R mod x^L = 1 (*this[0] ==
    1)
  for (int level = 0; (1 << level) < L; ++level) </pre>
   Poly 0 = R.Mul(Poly(data(), min<int>(2 << level,
    size())));
   Poly Q(2 \ll level); Q[0] = 1;
   for (int j = (1 << level); j < (2 << level); ++j)</pre>
   Q[j] = modsub(mod, O[j]);
   R = R.Mul(Q).isz(4 << level);
  }
  return R.isz(L);
 static int LinearRecursion(const V &a, const V &c,
    int64_t n) { // a_n = \sum c_j a_(n-j)}
  const int k = (int)a.size();
  assert((int)c.size() == k + 1);
  Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
  fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
C[k] = 1;
  while (n) {
   if (n % 2) W = W.Mul(M).DivMod(C).second;
   n /= 2, M = M.Mul(M).DivMod(C).second;
  int ret = 0;
  fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
  return ret:
 }
#undef fi
using Poly_t = Poly<998244353, 3, 1 << 20>;
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
5.20 Quadratic residue
struct S {
 int MOD. w:
 int64_t x, y;
 S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
  : MOD(m), w(w_{-}), x(x_{-}), y(y_{-}) {}
 S operator*(const S &rhs) const {
```

- d[i].begin();

```
int w_{-} = w;
                                                                         pivot(i, s);
  if (w_{-} == -1) w_{-} = rhs.w;
  assert(w_! = -1 \text{ and } w_ == rhs.w);
  return { MOD, w_,
(x * rhs.x + y * rhs.y % MOD * w) % MOD,
                                                                       if (!phase(0)) return VD(n, inf);
                                                                      VD x(n);
   (x * rhs.y + y * rhs.x) % MOD };
                                                                      for (int i = 0; i < m; ++i)</pre>
                                                                       if (p[i] < n) x[p[i]] = d[i][n + 1];
};
                                                                       return x;
int get_root(int n, int P) {
 if (P == 2 or n == 0) return n;
if (qpow(n, (P - 1) / 2, P) != 1) return -1;
                                                                      5.22 Simplex Construction
  auto check = [&](int x) {
                                                                     Standard form: maximize \sum_{1 < i < n} c_i x_i such that for all 1 \le j \le m,
  return qpow(x, (P - 1) / 2, P); };
if (check(n) == P-1) return -1;
                                                                     \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j and x_i \geq 0 for all 1 \leq i \leq n.
  int64_t a; int w; mt19937 rnd(7122);
                                                                        1. In case of minimization, let c_i' = -c_i
  do { a = rnd() % P;
  w = ((a * a - n) % P + P) % P;
                                                                        2. \sum_{1 < i < n} A_{ji} x_i \ge b_j \to \sum_{1 < i < n} -A_{ji} x_i \le -b_j
  } while (check(w) != P - 1);
  return qpow(S(P, w, a, 1), (P + 1) / 2).x;
                                                                        3. \sum_{1 < i < n} A_{ji} x_i = b_j
                                                                              • \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
5.21 Simplex
                                                                              • \sum_{1 \le i \le n} A_{ji} x_i \ge b_j
namespace simplex {
// maximize c^Tx under Ax <= B
                                                                        4. If x_i has no lower bound, replace x_i with x_i - x_i^\prime
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
                                                                           Geometry
using VD = vector<double>
using VVD = vector<vector<double>>;
                                                                     6.1 Basic Geometry
const double eps = 1e-9;
                                                                     #define IM imag
const double inf = 1e+9;
                                                                     #define RE real
int n. m:
                                                                     using lld = int64_t;
VVD d:
                                                                     using llf = long double;
vector<int> p, q;
void pivot(int r, int s) {
                                                                     using PT = std::complex<1ld>;
                                                                     using PTF = std::complex<llf>
double inv = 1.0 / d[r][s];
                                                                     auto toPTF(PT p) { return PTF{RE(p), IM(p)}; }
for (int i = 0; i < m + 2; ++i)
                                                                     int sgn(11d x) \{ return (x > 0) - (x < 0); \}
  for (int j = 0; j < n + 2; ++j)
                                                                     1ld dot(PT a, PT b) { return RE(conj(a) * b); }
1ld cross(PT a, PT b) { return IM(conj(a) * b); }
   if (i != r && j != s)
    d[i][j] -= d[r][j] * d[i][s] * inv;
                                                                     int ori(PT a, PT b, PT c) {
for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
                                                                      return sgn(cross(b - a, c - a));
d[r][s] = inv; swap(p[r], q[s]);
                                                                     bool operator<(const PT &a, const PT &b) {</pre>
                                                                      return RE(a) != RE(b) ? RE(a) < RE(b) : IM(a) < IM(b);
bool phase(int z) {
int x = m + z;
                                                                     int quad(PT p) {
  return (IM(p) == 0) // use sgn for PTF
while (true) {
  int s = -1;
  for (int i = 0; i <= n; ++i) {
  if (!z && q[i] == -1) continue;</pre>
                                                                        ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
                                                                     int argCmp(PT a, PT b) {
   if (s == -1) \mid d[x][i] < d[x][s]) s = i;
                                                                       // -1 / 0 / 1 <-> < / == / > (atan2)
                                                                       int qa = quad(a), qb = quad(b);
  if (d[x][s] > -eps) return true;
                                                                       if (qa != qb) return sgn(qa - qb);
                                                                       return sgn(cross(b, a));
  for (int i = 0; i < m; ++i) {</pre>
   if (d[i][s] < eps) continue;
if (r == -1 || \</pre>
                                                                      template <typename V> 11f area(const V & pt) {
                                                                      11d ret = 0;
    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
                                                                      for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
                                                                        ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
  if (r == -1) return false;
                                                                       return ret / 2.0;
  pivot(r, s);
                                                                     PT rot90(PT p) { return PT{-IM(p), RE(p)}; }
                                                                     PTF project(PTF p, PTF q) { // p onto q return dot(p, q) * q / dot(q, q);
VD solve(const VVD &a, const VD &b, const VD &c) {
m = b.size(), n = c.size();
 d = VVD(m + 2, VD(n + 2));
                                                                     11f FMOD(11f x) {
for (int i = 0; i < m; ++i)
                                                                       if (x < -PI) x += PI * 2;
 for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
                                                                      if (x > PI) x -= PI * 2;
p.resize(m), q.resize(n + 1);
for (int i = 0; i < m; ++i)
                                                                       return x;
  p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];
q[n] = -1, d[m + 1][n] = 1;
                                                                      6.2 Segment & Line Intersection
 int r = 0;
                                                                     struct Segment {
for (int i = 1; i < m; ++i)
if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
                                                                       PT st, dir; // represent st + t*dir for 0<=t<=1
                                                                       Segment(PT s, PT e) : st(s), dir(e - s) {}
 if (d[r][n + 1] < -eps) {</pre>
                                                                       static bool valid(lld p, lld q) {
                                                                        // is there t s.t. 0 <= t <= 1 && qt == p ?
  pivot(r, n);
  if (!phase(1) \mid \mid d[m + 1][n + 1] < -eps)
                                                                        if (q < 0) q = -q, p = -p;
   return VD(n, -inf);
                                                                        return 0 <= p && p <= q;
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
                                                                       }
   int s = min_element(d[i].begin(), d[i].end() - 1)
```

bool isInter(Segment A, PT P) {

```
if (A.dir == PT(0)) return P == A.st;
                                                                      flag[f.c][f.a] != flag[f.a][f.c])
 return cross(P - A.st, A.dir) == 0 &&
Segment::valid(dot(P - A.st, A.dir), norm(A.dir));
                                                                    next.emplace_back(f.c,f.a,i);
                                                                  now=next;
template <typename U, typename V>
                                                                 }
bool isInter(U A, V B) {
  if (cross(A.dir, B.dir) == 0)
                                                                 return now;
  return // handle parallel yourself
                                                                6.5 2D Farthest Pair
   isInter(A, B.st) || isInter(A, B.st+B.dir) ||
 isInter(B, A.st) || isInter(B, A.st+A.dir);
PT D = B.st - A.st;
                                                                // stk is from convex hull
                                                                n = (int)(stk.size());
 11d C = cross(A.dir, B.dir);
                                                                int pos = 1, ans = 0; stk.push_back(stk[0]);
                                                                for(int i=0;i<n;i++) {</pre>
 return U::valid(cross(D, A.dir), C) &&
                                                                 while(abs(cross(stk[i+1]-stk[i],
   V::valid(cross(D, B.dir), C);
                                                                   stk[(pos+1)%n]-stk[i])) >
struct Line
                                                                   abs(cross(stk[i+1]-stk[i],
PT st, ed, dir;
Line (PT s, PT e)
                                                                   stk[pos]-stk[i]))) pos = (pos+1)%n;
                                                                 ans = max({ans, dis(stk[i], stk[pos]),
  : st(s), ed(e), dir(e - s) {}
                                                                  dis(stk[i+1], stk[pos])});
PTF intersect(const Line &A, const Line &B) {
                                                                6.6 kD Closest Pair (3D ver.)
11f t = cross(B.st - A.st, B.dir) /
  llf(cross(A.dir, B.dir));
                                                                11f solve(vector<P> v) {
 return toPTF(A.st) + PTF(t) * toPTF(A.dir);
                                                                 shuffle(v.begin(), v.end(), mt19937());
                                                                 unordered_map<lld, unordered_map<lld,
                                                                  unordered_map<lld, int>>> m;
6.3 2D Convex Hull
                                                                 llf d = dis(v[0], v[1]);
// returns a convex hull in counterclockwise order
                                                                 auto Idx = [&d] (11f x) -> 11d {
                                                                  return round(x * 2 / d) + 0.1; }
// for a non-strict one, change cross >= to >
vector<PT> convex_hull(vector<PT> p) {
                                                                 auto rebuild_m = [&m, &v, &Idx](int k) {
 sort(all(p));
                                                                  m.clear();
                                                                  for (int i = 0; i < k; ++i)
m[Idx(v[i].x)][Idx(v[i].y)]</pre>
 if (p[0] == p.back()) return {p[0]};
 int n = p.size(), t = 0;
                                                                    [Idx(v[i].z)] = i
 vector<PT> h(n + 1);
 for (int _ = 2, s = 0; _--; s = --t, reverse(all(p)))
                                                                 }; rebuild_m(2);
  for (PT i : p) {
                                                                 for (size_t i = 2; i < v.size(); ++i) {</pre>
                                                                  const 11d kx = Idx(v[i].x), ky = Idx(v[i].y),
   while (t > s + 1 \&\& cross(i, h[t-1], h[t-2]) >= 0)
    t--
                                                                     kz = Idx(v[i].z); bool found = false;
                                                                  for (int dx = -2; dx <= 2; ++dx) {
   h[t++] = i;
                                                                   const 11d nx = dx + kx;
 return h.resize(t), h;
                                                                    if (m.find(nx) == m.end()) continue;
                                                                   auto& mm = m[nx];
                                                                   for (int dy = -2; dy <= 2; ++dy) {
6.4 3D Convex Hull
                                                                    const 11d ny = dy + ky;
// return the faces with pt indexes
                                                                    if (mm.find(ny) == mm.end()) continue;
int flag[MXN][MXN];
                                                                     auto& mmm = mm[ny];
                                                                    for (int dz = -2; dz <= 2; ++dz) {
struct Point{
 ld x, y, z;
                                                                     const 11d nz = dz + kz;
 Point operator * (const ld &b) const {
                                                                     if (mmm.find(nz) == mmm.end()) continue;
  return (Point) {x*b,y*b,z*b}; }
                                                                     const int p = mmm[nz];
 Point operator * (const Point &b) const {
                                                                     if (dis(v[p], v[i]) < d) {</pre>
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
                                                                       d = dis(v[p], v[i]);
                                                                       found = true;
Point ver(Point a, Point b, Point c) {
  return (b - a) * (c - a);}
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0;
                                                                  if (found) rebuild_m(i + 1);
 REP(i,n) REP(j,n) flag[i][j] = 0;
                                                                  else m[kx][ky][kz] = i;
 vector<Face> now;
 now.emplace_back(0,1,2);
                                                                 return d;
                                                                }
 now.emplace_back(2,1,0);
 for (int i=3; i<n; i++){</pre>
                                                                6.7 Simulated Annealing
 ftop++; vector<Face> next;
REP(j, SZ(now)) {
Face& f=now[j]; int ff = 0;
                                                                11f anneal() {
                                                                 mt19937 rnd_engine( seed );
   ld d=(pt[i]-pt[f.a]).dot(
                                                                 uniform_real_distribution< llf > rnd( 0, 1 );
     ver(pt[f.a], pt[f.b], pt[f.c]));
                                                                 const 11f dT = 0.001;
   if (d <= 0) next.push_back(f);</pre>
                                                                 // Argument p
   if (d > 0) ff=ftop;
                                                                 11f S_cur = calc( p ), S_best = S_cur;
for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
   else if (d < 0) ff=-ftop;</pre>
   flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
                                                                  // \dot{M}odify p to p_prime
                                                                  const 11f S_prime = calc( p_prime );
  REP(j, SZ(now)) {
  Face& f=now[j];
                                                                  const llf delta_c = S_prime - S_cur;
llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
   if (flag[f.a][f.b] > 0 &&
                                                                  if ( rnd( rnd_engine ) <= prob )</pre>
     flag[f.a][f.b] != flag[f.b][f.a])
                                                                   S_cur = S_prime, p = p_prime;
    next.emplace_back(f.a,f.b,i);
                                                                  if ( S_prime < S_best ) // find min</pre>
   if (flag[f.b][f.c] > 0 &&
                                                                   S_best = S_prime, p_best = p_prime;
     flag[f.b][f.c] != flag[f.c][f.b])
    next.emplace_back(f.b,f.c,i);
                                                                 return S_best;
   if (flag[f.c][f.a] > 0 &&
```

# 6.8 Half Plane Intersection

```
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
bool operator<(const Line &lhs, const Line &rhs) {</pre>
  if (int cmp = argCmp(lhs.dir, rhs.dir))
    return cmp == -1;
  return ori(lhs.st, lhs.ed, rhs.st) < 0;
// intersect function is in "Segment Intersect"
11f HPI(vector<Line> &lines) {
  sort(lines.begin(), lines.end());
  deque<Line> que;
  deque<PTF> pt;
  que.push_back(lines[0]);
  for (int i = 1; i < (int)lines.size(); i++) {</pre>
    if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
     continue;
#define POP(L, R) \
    while (pt.size() > 0 \
      && ori(L.st, L.ed, pt.back()) < 0) \
pt.pop_back(), que.pop_back(); \
    while (pt.size() > 0 \
      && ori(R.st, R.ed, pt.front()) < 0) \
pt.pop_front(), que.pop_front();
    POP(lines[i], lines[i]);
    pt.push_back(intersect(que.back(), lines[i]));
    que.push_back(lines[i]);
  POP(que.front(), que.back())
  if (que.size() <= 1 ||</pre>
    argCmp(que.front().dir, que.back().dir) == 0)
  pt.push_back(intersect(que.front(), que.back()));
  return area(pt);
```

### 6.9 Minkowski Sum

```
vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
hull(A), hull(B);
vector<pll> C(1, A[0] + B[0]), s1, s2;
for(int i = 0; i < SZ(A); ++i)
s1.pb(A[(i + 1) % SZ(A)] - A[i]);
for(int i = 0; i < SZ(B); i++)</pre>
 s2.pb(B[(i + 1) % SZ(B)] - B[i]);
 for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
 if (p2 >= SZ(B)
    | | (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
   C.pb(C.back() + s1[p1++]);
   C.pb(C.back() + s2[p2++]);
 return hull(C), C;
```

# 6.10 Circle Class

```
struct Circle { PTF o; llf r; };
vector<llf> intersectAngle(Circle A, Circle B) {
PTF dir = B.o - A.o; llf d2 = norm(dir);
if (norm(A.r - B.r) >= d2) // norm(x) := |x|^2
 if (A.r < B.r) return {-PI, PI}; // A in B</pre>
  else return {}; // B in A
if (norm(A.r + B.r) <= d2) return {};
llf dis = abs(dir), theta = arg(dir);</pre>
llf phi = acos((A.r * A.r + d2 - B.r * B.r) /
   (2 * A.r * dis));
11f L = FMOD(theta - phi), R = FMOD(theta + phi);
return { L, R };
vector<PTF> intersectPoint(Circle a, Circle b) {
llf d = abs(a.o - b.o);
 if (d >= b.r+a.r || d <= abs(b.r-a.r)) return {};</pre>
11f dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
PTF dir = (a.o - b.o) / d;
PTF u = dir*d1 + b.o;
PTF v = rot90(dir) * sqrt(max<llf>(0, b.r*b.r-d1*d1));
 return {u + v, u - v};
```

### 6.11 Intersection of line and Circle

```
vector<PTF> line_interCircle(const PTF &p1,
  const PTF &p2, const PTF &c, const double r) {
 PTF ft = p1 + project(c-p1, p2-p1), vec = p2-p1;
 llf dis = abs(c - ft);
 if (abs(dis - r) < eps) return {ft};</pre>
 if (dis > r) return {};
 vec = vec * sqrt(r * r - dis * dis) / abs(vec);
 return {ft + vec, ft - vec};
}
```

# 6.12 Intersection of Polygon and Circle

```
// Divides into multiple triangle, and sum up
// test by HDU2892
11f _area(PTF pa, PTF pb, llf r) {
 if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
 if (abs(pb) < eps) return 0;</pre>
 11f S, h, theta;
 11f a = abs(pb), b = abs(pa), c = abs(pb - pa);
 llf cosB = dot(pb, pb - pa) / a / c, B = acos(cosB); llf cosC = dot(pa, pb) / a / b, C = acos(cosC);
 if (a > r) {
  S = (C / 2) * r * r;
h = a * b * sin(C) / c;
  if (h < r && B < PI / 2)
   S = (acos(h / r) * r * r - h * sqrt(r*r - h*h));
   else if (b > r) {
  theta = PI - B - asin(sin(B) / r * a);
  S = 0.5 * a*r * sin(theta) + (C - theta) / 2 * r*r;
 } else
  S = 0.5 * sin(C) * a * b;
 return S;
11f area_poly_circle(const vector<PTF> &poly,
  const PTF &0, const llf r) {
 11f S = 0;
 for (int i = 0, N = poly.size(); i < N; ++i)</pre>
  S += _area(poly[i] - 0, poly[(i + 1) % N] - 0, r) *
     ori(0, poly[i], poly[(i + 1) % N]);
 return fabs(S);
```

# 6.13 Point & Hulls Tangent

} else {

```
#define above(P, Vi, Vj) (ori(P, Vi, Vj) > \theta) // true
    if Vi is above Vj
#define below(P, Vi, Vj) (ori(P, Vi, Vj) < 0) // true</pre>
    if Vi is below Vj
// Rtangent_PointPolyC(): binary search for convex
    polygon right tangent
    Input: P = a 2D point (exterior to the polygon)
//
        n = number of polygon vertices
        V = array of vertices for a 2D convex polygon
11
    with V[n] = V[0]
   Return: index "i" of rightmost tangent point V[i]
int Rtangent_PointPolyC(PT P, int n, PT *V) {
int a, b, c;
int upA, dnC;
 if (below(P, V[1], V[0]) && !above(P, V[n - 1], V[0]))
  return 0;
 for (a = 0, b = n;;) {
  c = (a + b) / 2;
  dnC = below(P, V[c + 1], V[c]);
  if (dnC && !above(P, V[c - 1], V[c]))
   return c;
  upA = above(P, V[a + 1], V[a]);
  if (upA) {
   if (dnC) {
    b = c;
   } else {
    if (above(P, V[a], V[c]))
     b = c;
    else
     a = c;
  } else {
   if (!dnC) {
```

while (ori(V[ix1], W[ix2], W[ix2 - 1]) >= 0) {

```
if (below(P, V[a], V[c]))
                                                                         // get Ltangent from V[ix1] to W
    b = c;
                                                               done = false; // not done if had to adjust this
   else
                                                              }
    a = c;
  }
                                                             *t1 = ix1;
 }
                                                             *t2 = ix2;
                                                             return;
}
                                                                  Tangent line of Two Circle
                                                            6.15
// Ltangent_PointPolyC(): binary search for convex
    polygon left tangent
                                                            vector<Line>
   Input: P = a 2D point (exterior to the polygon)
                                                            tanline(const Circle &c1, const Circle &c2, int sign1){
        n = number of polygon vertices
//
                                                             // sign1 = 1 for outer tang, -1 for inter tang
//
        V = array of vertices for a 2D convex polygon
                                                             vector<Line> ret;
    with V[n]=V[0]
                                                             if (norm(c1.o - c2.o) < eps) return ret;
// Return: index "i" of leftmost tangent point V[i]
                                                             11f d = abs(c1.o - c2.o);
int Ltangent_PointPolyC(PT P, int n, PT *V) {
                                                             PTF v = (c2.o - c1.o) / d;
int a, b, c;
                                                             11f c = (c1.r - sign1 * c2.r) / d;
int dnA, dnC;
                                                             if (c * c > 1) return ret;
                                                             llf h = sqrt(max<llf>(0, 1 - c * c));
if (above(P, V[n - 1], V[0]) && !below(P, V[1], V[0]))
                                                             for (int sign2 : {1, -1}) {
 return 0;
                                                              PTF n = c * v + sign2 * h * rot90(v);
                                                              PTF p1 = c1.o + n * c1.r;
for (a = 0, b = n;;) {
                                                              PTF p2 = c2.o + n * (c2.r * sign1);
 c = (a + b) / 2;
dnC = below(P, V[c + 1], V[c]);
                                                              if (norm(p2 - p1) < eps)
                                                               p2 = p1 + rot90(c2.o - c1.o);
 if (above(P, V[c - 1], V[c]) && !dnC)
                                                              ret.push_back({p1, p2});
 dnA = below(P, V[a + 1], V[a]);
                                                             return ret;
  if (dnA) {
  if (!dnC) {
   b = c;
                                                            6.16
                                                                  Minimum Covering Circle
   } else {
                                                            template<typename P>
   if (below(P, V[a], V[c]))
                                                            Circle getCircum(const P &a, const P &b, const P &c){
    b = c;
                                                             Real a1 = a.x-b.x, b1 = a.y-b.y;
    else
                                                             Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
     a = c;
                                                             Real a2 = a.x-c.x, b2 = a.y-c.y;
  }
                                                             Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
  } else {
                                                             Circle cc;
  if (dnC) {
                                                             cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
   a = c;
                                                             cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
   } else {
                                                             cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
   if (above(P, V[a], V[c]))
                                                             return cc;
    b = c;
   else
     a = c;
                                                            template<typename P>
                                                            Circle MinCircleCover(const vector<P>& pts){
                                                             random_shuffle(pts.begin(), pts.end());
                                                             Circle c = \{ pts[0], 0 \};
                                                             for(int i=0;i<(int)pts.size();i++){</pre>
6.14
      Convex Hulls Tangent
                                                              if (dist(pts[i], c.o) <= c.r) continue;</pre>
                                                              c = { pts[i], 0 };
for (int j = 0; j < i; j++) {
// RLtangent_PolyPolyC(): get the RL tangent between
    two convex polygons
                                                               if(dist(pts[j], c.o) <= c.r) continue;</pre>
   Input: m = number of vertices in polygon 1
                                                               c.o = (pts[i] + pts[j]) / 2;
        V = array of vertices for convex polygon 1 with
//
                                                               c.r = dist(pts[i], c.o)
     V[m]=V[0]
                                                               for (int k = 0; k < j; k++) {
11
        n = number of vertices in polygon 2
                                                                if (dist(pts[k], c.o) <= c.r) continue;</pre>
11
        W = array of vertices for convex polygon 2 with
                                                                c = getCircum(pts[i], pts[j], pts[k]);
     W[n]=W[0]
   Output: *t1 = index of tangent point V[t1] for
                                                              }
    polygon 1
        *t2 = index of tangent point W[t2] for polygon
                                                             return c;
void RLtangent_PolyPolyC(int m, PT *V, int n, PT *W,
    int *t1, int *t2) {
                                                                  KDTree (Nearest Point)
int ix1, ix2; // search indices for polygons 1 and 2
                                                            const int MXN = 100005;
                                                            struct KDTree {
// first get the initial vertex on each polygon
ix1 = Rtangent_PointPolyC(W[0], m, V); // right
                                                             struct Node {
    tangent from W[0] to V
                                                              int x,y,x1,y1,x2,y2;
ix2 = Ltangent_PointPolyC(V[ix1], n, W); // left
                                                              int id,f;
Node *L, *R;
    tangent from V[ix1] to W
                                                             } tree[MXN], *root;
// ping-pong linear search until it stabilizes
int done = false; // flag when done
                                                             LL dis2(int x1, int y1, int x2, int y2) {
while (done == false) {
                                                              LL dx = x1-x2, dy = y1-y2;
                                                              return dx*dx+dy*dy;
 done = true; // assume done until..
 while (ori(W[ix2], V[ix1], V[ix1 + 1]) <= 0) {</pre>
  ++ix1; // get Rtangent from W[ix2] to V
                                                             static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
                                                             static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
```

void init(vector<pair<int,int>> ip) {

```
n = ip.size();
                                                                 for (int i = 0; i < n; ++i) p[idx[i]] = i;
  for (int i=0; i<n; i++) {</pre>
                                                                 for (auto [i, j]: 1) {
                                                                  // do here
   tree[i].id = i;
   tree[i].x = ip[i].first;
                                                                  swap(p[i], p[j]);
   tree[i].y = ip[i].second;
                                                                  idx[p[i]] = i, idx[p[j]] = j;
 root = build_tree(0, n-1, 0);
                                                                6.19
                                                                       Circle Cover
Node* build_tree(int L, int R, int d) {
 if (L>R) return nullptr
                                                                const int N = 1021;
  int M = (L+R)/2; tree[M].f = d%2;
                                                                struct CircleCover {
  nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
                                                                 int C
  tree[M].x1 = tree[M].x2 = tree[M].x;
                                                                 Cir c[N]
  tree[M].y1 = tree[M].y2 = tree[M].y
                                                                 bool g[N][N], overlap[N][N];
  tree[M].L = build_tree(L, M-1, d+1);
                                                                 // Area[i] : area covered by at least i circles
  if (tree[M].L) {
                                                                 double Area[ N ];
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
                                                                 void init(int _C){ C = _C;}
   tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
                                                                 struct Teve
   tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
                                                                  PTF p; double ang; int add;
Teve() {}
   tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
                                                                  Teve(PTF _a, double _b, int _c):p(_a), ang(_b), add(
  tree[M].R = build_tree(M+1, R, d+1);
                                                                    _c){}
  if (tree[M].R) {
                                                                  bool operator<(const Teve &a)const
  tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
                                                                  {return ang < a.ang;}
   tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
                                                                 }eve[N * 2];
  tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
                                                                 // strict: x = 0, otherwise x = -1
                                                                 bool disjuct(Cir &a, Cir &b, int x)
                                                                 {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
bool contain(Cir &a, Cir &b, int x)
 return tree+M:
                                                                 {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
int touch(Node* r, int x, int y, LL d2){
                                                                 bool contain(int i, int j) {
 LL dis = sqrt(d2)+1;
                                                                  /* c[j] is non-strictly in c[i]. */
  if (x<r->x1-dis || x>r->x2+dis ||
                                                                  return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c[i].R) \mid c[i].R - c[i].R)
    y<r->y1-dis || y>r->y2+dis)
                                                                    [j].R) == 0 \&\& i < j)) \&\& contain(c[i], c[j], -1);
   return 0;
  return 1;
                                                                 void solve(){
                                                                  fill_n(Area, C + 2, 0);
 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
                                                                  for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)
  if (!r || !touch(r, x, y, md2)) return;
 LL d2 = dis2(r->x, r->y, x, y)
                                                                    overlap[i][j] = contain(i, j);
  if (d2 < md2 || (d2 == md2 && mID < r->id)) {
                                                                  for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)
  mID = r -> id;
  md2 = d2;
                                                                    g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                                                                       disjuct(c[i], c[j], -1));
  // search order depends on split dim
                                                                  for(int i = 0; i < C; ++i){</pre>
  if ((r->f == 0 && x < r->x) ||
                                                                   int E = 0, cnt = 1;
    (r->f == 1 \&\& y < r->y)) {
                                                                   for(int j = 0; j < C; ++j)
   nearest(r->L, x, y, mID, md2);
                                                                    if(j != i && overlap[j][i])
   nearest(r->R, x, y, mID, md2);
                                                                     ++cnt:
  } else {
                                                                   for(int j = 0; j < C; ++j)
  nearest(r->R, x, y, mID, md2);
                                                                    if(i != j && g[i][j]) {
   nearest(r->L, x, y, mID, md2);
                                                                     auto IP = intersectPoint(c[i], c[j]);
                                                                     PTF aa = IP[0], bb = IP[1];
                                                                     llf A = arg(aa-c[i].0), B = arg(bb-c[i].0);
eve[E++] = Teve(bb,B,1), eve[E++]=Teve(aa,A,-1);
int query(int x, int y) {
 int id = 1029384756:
                                                                     if(B > A) ++cnt;
 LL d2 = 102938475612345678LL;
 nearest(root, x, y, id, d2);
                                                                   if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
  return id;
                                                                    sort(eve, eve + E);
} tree;
                                                                    eve[E] = eve[0];
                                                                    for(int j = 0; j < E; ++j){
6.18
      Rotating Sweep Line
                                                                     cnt += eve[j].add;
void rotatingSweepLine(pair<int, int> a[], int n) {
                                                                     Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
 vector<pair<int, int>> 1;
                                                                     double theta = eve[j + 1].ang - eve[j].ang;
1.reserve(n * (n - 1) / 2);
                                                                     if (theta < 0) theta += 2. * pi;</pre>
for (int i = 0; i < n; ++i)
                                                                     Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
  for (int j = i + 1; j < n; ++j)
   1.emplace_back(i, j);
 sort(1.begin(), 1.end(), [&a](auto &u, auto &v){
 1ld udx = a[u.first].first - a[u.second].first;
1ld udy = a[u.first].second - a[u.second].second;
 11d vdx = a[v.first].first - a[v.second].first;
  11d vdy = a[v.first].second - a[v.second].second;
                                                                     Stringology
  if (udx == 0 or vdx == 0) return not udx == 0;
  int s = sgn(udx * vdx);
                                                                7.1 Suffix Array
  return udy * vdx * s < vdy * udx * s;
 });
                                                                namespace sfx {
vector<int> idx(n), p(n);
                                                                bool _t[maxn * 2];
iota(idx.begin(), idx.end(), 0);
sort(idx.begin(), idx.end(), [&a](int i, int j){
                                                                int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], _c[maxn * 2];
                                                                int x[maxn], _p[maxn], _q[maxn * 2];
 return a[i] < a[j]; });
```

```
// sa[i]: sa[i]-th suffix is the
                                                                   if (!last) {
                                                                     st[cur].fail = root;
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix
                                                                   } else {
// of suffix sa[i] and suffix sa[i - 1]
                                                                     int q = st[last].ch[c];
void pre(int *a, int *c, int n, int z) {
                                                                     if (st[q].len == st[last].len + 1) {
memset(a, 0, sizeof(int) * n);
                                                                       st[cur].fail = q;
memcpy(x, c, sizeof(int) * z);
                                                                     } else {
                                                                       int clone = ++tot;
void induce(int *a,int *c,int *s,bool *t,int n,int z){
                                                                       st[clone] = st[q];
memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
                                                                       st[clone].len = st[last].len + 1;
                                                                       st[st[cur].fail = st[q].fail = clone].cnt = 0;
 if (a[i] && !t[a[i] - 1])
                                                                       while (last && st[last].ch[c] == q) {
                                                                          st[last].ch[c] = clone;
   a[x[s[a[i] - 1]]++] = a[i] - 1;
memcpy(x, c, sizeof(int) * z);
                                                                          last = st[last].fail;
for (int i = n - 1; i >= 0; --i)
if (a[i] && t[a[i] - 1])
                                                                     }
   a[--x[s[a[i] - 1]]] = a[i] - 1;
                                                                   st[last = cur].cnt += 1;
void sais(int *s, int *a, int *p, int *q,
bool *t, int *c, int n, int z) {
                                                                  void init(const char* s) {
bool uniq = t[n - 1] = true;
                                                                   root = last = tot = 1;
int nn=0, nmxz=-1, *nsa = a+n, *ns=s+n, last=-1;
                                                                   st[root] = node(0);
memset(c, 0, sizeof(int) * z);
                                                                   for (char c; c = *s; ++s) extend(c - 'a');
for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
 for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
                                                                  int q[N];
if (uniq) {
                                                                  void dp() {
                                                                   for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
 for (int i = 0; i < n; ++i) a[--c[s[i]]] = i;
  return;
                                                                   int head = 0, tail = 0;
                                                                   for (int i = 1; i <= tot; i++)</pre>
 for (int i = n - 2; i \ge 0; --i)
 t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
                                                                     if (st[i].indeg == 0) q[tail++] = i;
pre(a, c, n, z);
for (int i = 1; i <= n - 1; ++i)
                                                                   while (head != tail) {
                                                                     int now = q[head++]
 if (t[i] && !t[i - 1])
                                                                     if (int f = st[now].fail) {
   a[--x[s[i]]] = p[q[i] = nn++] = i;
                                                                       st[f].cnt += st[now].cnt;
induce(a, c, s, t, n, z);
for (int i = 0; i < n; ++i)
                                                                       if (--st[f].indeg == 0) q[tail++] = f;
  if (a[i] && t[a[i]] && !t[a[i] - 1]) {
                                                                   }
 bool neq = last < 0 || \</pre>
  memcmp(s + a[i], s + last,
(p[q[a[i]] + 1] - a[i]) * sizeof(int));
                                                                  int run(const char* s) {
                                                                   int now = root;
 ns[q[last = a[i]]] = nmxz += neq;
                                                                   for (char c; c = *s; ++s)
                                                                     if (!st[now].ch[c -= 'a']) return 0;
}}
sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
                                                                     now = st[now].ch[c];
pre(a, c, n, z);
for (int i = nn - 1; i >= 0; --i)
                                                                   return st[now].cnt;
  a[--x[s[p[nsa[i]]]]] = p[nsa[i]];
                                                                } SAM;
induce(a, c, s, t, n, z);
                                                                 7.3 Z value
void build(const string &s) {
const int n = int(s.size());
                                                                 vector<int> Zalgo(const string &s) {
for (int i = 0; i < n; ++i) _s[i] = s[i];
                                                                  vector<int> z(s.size(), s.size());
 _s[n] = 0; // s shouldn't contain 0
                                                                  for (int i = 1, 1 = 0, r = 0; i < z[0]; ++i) {
  int j = clamp(r - i, 0, z[i - 1]);</pre>
sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
for(int i = 0; i < n; ++i) rev[sa[i] = sa[i+1]] = i;</pre>
                                                                   for (; i + j < z[0] \text{ and } s[i + j] == s[j]; ++j);
int ind = hi[0] = 0:
                                                                   if (i + (z[i] = j) > r) r = i + z[1 = i];
for (int i = 0; i < n; ++i) {
 if (!rev[i]) {
                                                                  return z;
   ind = 0:
                                                                }
   continue;
                                                                 7.4 Manacher
 while (i + ind < n && \</pre>
                                                                int z[maxn];
   s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
                                                                int manacher(const string& s) {
  string t = ".";
 hi[rev[i]] = ind ? ind-- : 0;
                                                                  for(char c: s) t += c, t += '.';
int 1 = 0, r = 0, ans = 0;
                                                                  for (int i = 1; i < t.length(); ++i) {
z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
7.2 Suffix Automaton
struct SuffixAutomaton {
                                                                   while (i - z[i] \ge 0 \&\& i + z[i] < t.length()) {
struct node {
                                                                    if(t[i - z[i]] == t[i + z[i]]) ++z[i];
 int ch[K], len, fail, cnt, indeg;
node(int L = 0) : ch{}, len(L), fail(0), cnt(0),
                                                                    else break:
    indeg(0) {}
                                                                   if (i + z[i] > r) r = i + z[i], l = i;
 } st[N];
 int root, last, tot;
                                                                  for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
void extend(int c) {
                                                                  return ans;
 int cur = ++tot;
  st[cur] = node(st[last].len + 1);
 while (last && !st[last].ch[c]) {
                                                                 7.5 Lexico Smallest Rotation
    st[last].ch[c] = cur;
    last = st[last].fail;
                                                                string mcp(string s) {
```

int n = s.length();

```
s += s; int i = 0, j = 1;
while (i < n && j < n) {
  int k = 0;
  while (k < n && s[i + k] == s[j + k]) k++;
  ((s[i + k] <= s[j + k]) ? j : i) += k + 1;
  j += (i == j);
}
return s.substr(i < n ? i : j, n);
}</pre>
```

### 7.6 BWT

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a
vector<int> v[ SIGMA ];
void BWT(char* ori, char* res){
  // make ori -> ori + ori
  // then build suffix array
 }
 void iBWT(char* ori, char* res){
  for( int i = 0 ; i < SIGMA ; i ++ )</pre>
   v[ i ].clear();
  int len = strlen( ori );
for( int i = 0 ; i < len ; i ++ )</pre>
   v[ ori[i] - BASE ].push_back( i );
  vector<int> a:
                     ptr = 0 ; i < SIGMA ; i ++ )
  for(int i = 0,
   for( auto j : v[ i ] ){
    a.push_back( j );
    ori[ ptr ++ ] = BASE + i;
  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
   ptr = a[ ptr ];
  res[ len ] = 0;
}
} bwt;
```

# 7.7 Palindromic Tree

```
struct palindromic_tree{
struct node{
  int next[26],f,len;
  int cnt, num, st, ed;
 node( \texttt{int} \ l=\emptyset) : f(\emptyset), len(1), cnt(\emptyset), num(\emptyset) \ \{
  memset(next, 0, sizeof(next)); }
vector<node> st;
vector<char> s;
int last,n;
void init(){
 st.clear();s.clear();last=1; n=0;
 st.push_back(0);st.push_back(-1);
  st[0].f=1;s.push_back(-1); }
 int getFail(int x){
 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
  return x;}
void add(int c){
  s.push_back(c-='a'); ++n;
  int cur=getFail(last);
  if(!st[cur].next[c]){
  int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
   st[now].num=st[st[now].f].num+1;
  last=st[cur].next[c];
  ++st[last].cnt;}
 void dpcnt() {
  for (int i=st.size()-1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
int size(){ return st.size()-2;}
} pt;
int main() {
string s; cin >> s; pt.init();
for (int i=0; i<SZ(s); i++)</pre>
 int prvsz = pt.size(); pt.add(s[i]);
  if (prvsz != pt.size()) {
   int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [l,r]: s.substr(l, r-l+1)
```

```
}
}
return 0;
}
```

# 8 Misc

#### 8.1 Theorems

#### 8.1.1 Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii}=d(i)$ ,  $L_{ij}=-c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .

#### 8.1.2 Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on C

#### 8.1.3 Cayley's Formula

- Given a degree sequence  $d_1,d_2,\dots,d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .

#### 8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all  $1 \le k \le n$ .

#### 8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

#### 8.1.6 Euler's planar graph formula

```
V - E + F = C + 1, E \le 3V - 6(?)
```

#### 8.1.7 Pick's theorem

For simple polygon, when points are all integer, we have  $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$ 

#### 8.1.8 Matroid Intersection

Given matroids  $M_1=(G,I_1),M_2=(G,I_2)$ , find maximum  $S\in I_1\cap I_2$ . For each iteration, build the directed graph and find a shortest path from s to t.

- $s \to x : S \sqcup \{x\} \in I_1$
- $x \to t : S \sqcup \{x\} \in I_2$
- $y \to x : S \setminus \{y\} \sqcup \{x\} \in I_1$  (y is in the unique circuit of  $S \sqcup \{x\}$ )
- $x \to y : S \setminus \{y\} \sqcup \{x\} \in I_2$  (y is in the unique circuit of  $S \sqcup \{x\}$ )

Alternate the path, and |S| will increase by 1. Let  $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|$ . In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to  $x\in S$  and  $x\notin S$ , resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

#### 8.2 Convex 1D/1D DP

```
struct segment {
  int i, 1, r;
  segment() {}
  segment(int a, int b, int c): i(a), l(b), r(c) {}
};
  inline lld f(int l, int r){return dp[l] + w(l+1, r);}
  void solve() {
    dp[0] = 0;
    deque<segment> dq; dq.push_back(segment(0, 1, n));
    for (int i = 1; i <= n; ++i) {
        dp[i] = f(dq.front().i, i);
        while(dq.size()&&dq.front().r<i+1) dq.pop_front();
        dq.front().l = i + 1;
        segment seg = segment(i, i + 1, n);</pre>
```

```
while (dq.size() &&
                                                                   temp_v=par[temp_v];
   f(i, dq.back().1) < f(dq.back().i, dq.back().1))
                                                                  g[bcc_id+n].push_back(v);
    dq.pop_back();
                                                                  g[v].push_back(bcc_id+n);
  if (dq.size()) {
   int d = 1 << 20, c = dq.back().1;</pre>
                                                                  reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
   while (d \gg 1) if (c + d \ll dq.back().r)
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
  if (seg.1 <= n) dq.push_back(seg);</pre>
                                                             int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
                                                             void dfs(int u,int fa){
                                                              if(u<=n){
                                                                for(int i=0;i<(int)g[u].size();i++){</pre>
     ConvexHull Optimization
8.3
                                                                 int v=g[u][i];
                                                                 if(v==fa) continue;
struct L
                                                                 dfs(v,u);
mutable int64_t a, b, p;
                                                                 memset(tp,0x8f,sizeof tp);
bool operator<(const L &r) const { return a < r.a; }</pre>
                                                                 if(v<=n){
bool operator<(int64_t x) const { return p < x; }</pre>
                                                                  tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
                                                                  tp[1]=max(
struct DynamicHull : multiset<L, less<>>> {
                                                                   dp[u][0]+dp[v][0]+1
static const int64_t kInf = 1e18;
                                                                   dp[u][1]+max(dp[v][0],dp[v][1])
bool Isect(iterator x, iterator y)
 auto Div = [](int64_t a, int64_t b) {
                                                                 }else{
    return a / b - ((a ^ b) < 0 && a % b); }
                                                                  tp[0]=dp[u][0]+dp[v][0];
  if (y == end()) { x->p = kInf; return false;
                                                                  tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
  if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
                                                                 dp[u][0]=tp[0],dp[u][1]=tp[1];
  return x->p >= y->p;
                                                                }
                                                               }else{
void Insert(int64_t a, int64_t b) {
                                                                for(int i=0;i<(int)g[u].size();i++){</pre>
 auto z = insert({a, b, 0}), y = z++, x = y;
while (Isect(y, z)) z = erase(z);
                                                                 int v=g[u][i];
                                                                 if(v==fa) continue;
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
                                                                 dfs(v,u);
 while ((y = x) != begin() && (--x)->p >= y->p)
  Isect(x, erase(y));
                                                                min_dp[0][0]=0;
                                                                min_dp[1][1]=1;
int64_t Query(int64_t x) {
                                                                min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
 auto 1 = *lower_bound(x);
                                                                for(int i=0;i<(int)g[u].size();i++){</pre>
  return 1.a * x + 1.b;
                                                                int v=g[u][i];
                                                                 if(v==fa) continue;
                                                                 memset(tmp,0x8f,sizeof tmp);
                                                                 tmp[0][0]=max(
8.4
      Josephus Problem
                                                                  \min_{dp[0][0]+\max(dp[v][0],dp[v][1])}
// n people kill m for each turn
                                                                  min_dp[0][1]+dp[v][0]
int f(int n, int m) {
int s = 0;
                                                                 tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
for (int i = 2; i <= n; i++)
                                                                 tmp[1][0]=max(
 s = (s + m) \% i;
                                                                  min_dp[1][0]+max(dp[v][0],dp[v][1]),
 return s;
                                                                  min_dp[1][1]+dp[v][0]
// died at kth
                                                                 tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
int kth(int n, int m, int k){
                                                                 memcpy(min_dp,tmp,sizeof tmp);
if (m == 1) return n-1;
for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
                                                                dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
return k;
                                                               dp[u][0]=min_dp[0][0];
8.5 Cactus Matching
                                                             int main(){
                                                              int m,a,b;
vector<int> init_g[maxn],g[maxn*2];
                                                               scanf("%d%d",&n,&m);
for(int i=0;i<m;i++){</pre>
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u)
                                                                scanf("%d%d",&a,&b);
dfn[u]=low[u]=++dfs_idx;
                                                                init_g[a].push_back(b);
for(int i=0;i<(int)init_g[u].size();i++){</pre>
  int v=init_g[u][i];
                                                                init_g[b].push_back(a);
  if(v==par[u]) continue;
                                                               par[1]=-1;
  if(!dfn[v]){
   par[v]=u;
                                                               tarjan(1);
                                                               dfs(1,-1);
   tarjan(v);
                                                               printf("%d\n", max(dp[1][0], dp[1][1]));
   low[u]=min(low[u],low[v]);
                                                               return 0;
   if(dfn[u]<low[v]){</pre>
    g[u].push_back(v);
    g[v].push_back(u);
                                                             8.6 Tree Knapsack
                                                             int dp[N][K]; PII obj[N];
  }else{
   low[u]=min(low[u],dfn[v]);
                                                             vector<int> G[N];
   if(dfn[v]<dfn[u]){</pre>
                                                             void dfs(int u, int mx){
for(int s: G[u]) {
    int temp_v=u;
    bcc_id++;
                                                                if(mx < obj[s].first) continue;</pre>
    while(temp_v!=v){
                                                                for(int i=0;i<=mx-obj[s].FF;i++)</pre>
     g[bcc_id+n].push_back(temp_v);
                                                                 dp[s][i] = dp[u][i];
                                                                dfs(s, mx - obj[s].first);
     g[temp_v].push_back(bcc_id+n);
```

```
for(int i=obj[s].FF;i<=mx;i++)</pre>
   dp[u][i] = max(dp[u][i],
    dp[s][i - obj[s].FF] + obj[s].SS);
8.7 N Queens Problem
vector< int > solve( int n ) {
// no solution when n=2, 3
 vector< int > ret;
if ( n % 6 == 2 ) {
  for ( int i = 2 ; i <= n ; i += 2 )
    ret.push_back( i );</pre>
  ret.push_back( 3 ); ret.push_back( 1 );
for ( int i = 7 ; i <= n ; i += 2 )
  ret.push_back( i );
  ret.push_back( 5 );
 } else if ( n % 6 == 3 ) {
 for ( int i = 4 ; i <= n ; i += 2 )
  ret.push_back( i );
  ret.push_back( 2 );
  for ( int i = 5 ; i <= n ; i += 2 )
  ret.push_back( i );
  ret.push_back( 1 ); ret.push_back( 3 );
 } else {
  for ( int i = 2 ; i <= n ; i += 2 )
  ret.push_back( i );
for ( int i = 1 ; i <= n ; i += 2 )
   ret.push_back( i );
 return ret;
8.8 Binary Search On Fraction
struct Q {
11 p, q;
Q go(Q b, 11 d) \{ return \{ p + b.p*d, q + b.q*d \}; \}
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
 assert(pred(hi));
 bool dir = 1, L = 1, H = 1;
 for (; L || H; dir = !dir) {
    11 len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)
   if (Q mid = hi.go(lo, len + step);
  mid.p > N || mid.q > N || dir ^ pred(mid))
    t++;
   else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
 return dir ? hi : lo;
```