Contents Basic 1.1 vimrc 1 Basic 5.29 Golden Ratio Search* 15 se is nu ru et tgc sc hls cin cino+=j1 sw=2 sts=2 bs=2 vimrc Debug Macro..... 1.1 15 6 Geometry mouse=a "encoding=utf-8 ls=2 1.2 syn on | colo desert | filetype indent on Basic Geometry.... 15 SVG Writer 1.3 inoremap {<CR> {<CR>}<ESC>0 2D Convex Hull 16 Pragma Optimization 1.4 1 map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -g -std=gnu++20 -2D Farthest Pair 6.3 16 1.5 IO Optimization.... DCKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -MinMax Enclosing 2 Data Structure Rect Wconversion -fsanitize=address,undefined,float-Minkowski Sum..... 16 divide-by-zero,float-cast-overflow && echo success< 2.1 Dark Magic 6.6 Segment Intersection 16 Link-Cut Tree Halfplane Intersecmap <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -O2 -g -std=gnu 2.3 LiChao Segtree tion 16 ++20 && echo success<CR> Treap* SegmentDist map <F10> <ESC>:!./"%<"<CR> Linear Basis* 25 (Sausage) ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space Binary Search on Rotating Sweep Line 17 :]' \| md5sum \| cut -c-6 Segtree 3 6.10 Hull Cut* 17 let c_no_curly_error=1 3 3 Graph 6.11 Point In Hull 17 " setxkbmap -option caps:ctrl_modifier SCC 2-SAT 6.12 Point In Polygon 17 1.2 Debug Macro [a45c59] 3.2 6.13 Point In Polygon #define all(x) begin(x), end(x) BCC 3.3 Round Square Tree.. #ifdef CKISEKI 3.4 6.14 Cyclic Ternary Search 18 #include <experimental/iterator> 3.5 Edge TCC 6.15 Tangent of Points #define safe cerr<<__PRETTY_FUNCTION__<<" line "<</pre> Bipolar Orientation.. 3.6 to Hull 18 __LINE__<<" safe\n" 3.7 DMST..... 6.16 Circle Class & Inter-Dominator Tree #define debug(a...) debug_(#a, a) 3.8 section #define orange(a...) orange_(#a, a) 3.9 Edge Coloring Circle Common Tangent void debug_(auto s, auto ...a) { cerr << "\e[1;32m(" << s << ") = (";</pre> 18 3.10 Centroid Decomp.*.. 6.18 Line-Circle Inter-3.11 Lowbit Decomp. 5 section int f = 0; 3.12 Virtual Tree* (..., (cerr << (f++ ? ", " : "") << a)); 6.19 Poly-Circle Inter-3.13 Tree Hashing section 18 cerr << ")\e[0m\n";</pre> 3.14 Mo's Algo on Tree ... 6.20 Min Covering Circle . 18 3.15 Count Cycles 6.21 Circle Union 18 void orange_(auto s, auto L, auto R) { cerr << "\e[1;33m[" << s << "] = [";</pre> 3.16 Maximal Clique 6.22 Polygon Union..... 19 3.17 Maximum Clique ... 6.23 3D Point using namespace experimental; 19 6.24 3D Convex Hull copy(L, R, make_ostream_joiner(cerr, ", ")); 3.18 Min Mean Cycle 6.25 3D Projection 19 cerr << "]\e[0m\n";</pre> 3.19 Eulerian Trail 6.26 3D Skew Line Near-4 Flow & Matching #else HopcroftKarp..... 19 4.1 #define safe ((void)0) 6.28 Build Voronoi 20 Kuhn Munkres* **#define** debug(...) safe 6.29 kd Tree (Nearest Point)*.... Flow Models..... #define orange(...) safe Dinic 4.4 #endif 4.5 HLPP 6.30 kd Closest Pair (3D 1.3 SVG Writer [85759e] Global Min-Cut GomoryHu Tree..... ver.)* 20 4.6 #ifdef CKISEKI 4.7 6.31 Simulated Annealing* 21 class SVG { MCMF 4.8 6.32 Triangle Centers* ... 21 void p(string_view s) { o << s; }</pre> Dijkstra Cost Flow . . . 7 Stringology 21 void p(string_view s, auto v, auto... vs) { 4.10 Min Cost Circulation. **auto** i = s.find('\$'); Hash 4.11 General Matching ... 10 o << s.substr(0, i) << v, p(s.substr(i + 1), vs...); Suffix Array 4.12 Weighted Matching . 10 Suffix Array Tools* . . . 7.3 21 5 Math 11 ofstream o; string c = "red"; 7.4 21 Common Bounds ... 5.1 public: 5.2 Equations 11 SVG(auto f,auto x1,auto y1,auto x2,auto y2) : o(f) { Integer Division* Manacher p("<svg xmlns='http://www.w3.org/2000/svg' "viewBox='\$ \$ \$'>\n" Lyndon Factorization 5.4 FloorSum 11 ModMin Main Lorentz* 5.5 "<style>*{stroke-width:0.5%;}</style>\n", 7.10 BWT*..... 5.6 Floor Monoid Product x1, -y2, x2 - x1, y2 - y1);} 5.7 7.11 Palindromic Tree* ... 22 ax+by=gcd ~SVG() { p("</svg>\n"); } void color(string nc) { c = nc; } 5.8 Chinese Remainder . 8 Misc 23 5.9 DiscreteLog void line(auto x1, auto y1, auto x2, auto y2) { p("<line x1='\$' y1='\$' x2='\$' y2='\$' stroke='\$'/>\n", 8.1 23 5.10 Quadratic Residue ... Stable Marriage..... 23 x1, -y1, x2, -y2, c); } FWT... 5.11 FWT..... 5.12 Packed FFT 8.3 Weight Matroid Invoid circle(auto x, auto y, auto r) { p("<circle cx='\$' cy='\$' r='\$' stroke='\$' "</pre> tersection* 5.13 CRT for arbitrary mod 12 Bitset LCS "fill='none'/>\n", x, -y, r, c); } Prefix Substring LCS. 24 5.14 NTT/FFT 8.5 5.15 Formal Power Series 8.6 Convex 1D/1D DP* ... 24 void text(auto x, auto y, string s, int w = 12) { p("<text x='\$' y='\$' font-size='\$px'>\$</text>\n", 5.16 Partition Number ... ConvexHull Opti-5.17 Pi Count 13 mization $x, -y, w, s); }$ 5.18 Min 25 Sieve* 14 Min Plus Convolution 24 }; // write wrapper for complex if use complex 5.19 Miller Rabin SMAWK..... #else 5.20 Pollard Rho 8.10 De-Bruijn 24 struct SVG { SVG(auto ...) {} }; // you know how to 5.21 Montgomery 8.11 Josephus Problem . . 24 #endif 5.22 Berlekamp Massev . . 14 1.4 Pragma Optimization [6006f6] 8.12 N Queens Problem* . 25 5.23 Gauss Elimination... 14 $\verb"pragma" GCC" optimize("Ofast, no-stack-protector")$ 8.13 Manhattan MST.... 25 5.24 CharPoly 8.14 Binary Search On #pragma GCC optimize("no-math-errno,unroll-loops") Fraction..... #pragma GCC target("sse,sse2,sse3,ssse3,sse4") 5.26 Simplex Construction 15 8.15 Cartesian Tree 25 #pragma GCC target("popcnt,abm,mmx,avx,arch=skylake") 5.27 Adaptive Simpson .. 15 8.16 Nim Product 25 _builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040) 8.17 Grid 5.28 Poly Roots*..... 15 1.5 IO Optimization [c9494b]

| static inline int gc() {

int find_root(int u) {

```
constexpr int B = 1<<20; static char buf[B], *p, *q;</pre>
                                                               int la = 0;
 if (p == q) q = (p = buf) + fread(buf, 1, B, stdin);
                                                               for (access(u); u; u = lc) down(la = u);
return q == buf ? EOF : *p++;
                                                               return la;
                                                              void split(int x, int y) { chroot(x); access(y); }
2
     Data Structure
                                                              void chroot(int u) { access(u); set_rev(u); }
   Dark Magic [095f25]
                                                              /* SPLIT_HASH_HERE */
#include <ext/pb_ds/assoc_container.hpp>
                                                             public:
#include <ext/pb_ds/priority_queue.hpp>
                                                              LCT(int n = 0) : o(n + 1) {}
using namespace __gnu_pbds;
                                                              void set_val(int u, const Val &v) {
  splay(++u); cur.v = v; up(u); }
// heap tags: paring/binary/binomial/rc_binomial/thin
template<typename T>
                                                              void set_sval(int u, const SVal &v) {
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>, \
                                                               access(++u); cur.sv = v; up(u); }
                  pairing_heap_tag>;
                                                              Val query(int x, int y) {
  ′pbds_heap::point_iterator
                                                               split(++x, ++y); return o[y].prod; }
// x = pq.push(10); pq.modify(x, 87); a.join(b);
                                                              SVal subtree(int p, int u) {
// tree tags: rb_tree_tag/ov_tree_tag/splay_tree_tag
                                                               chroot(++p); access(++u); return cur.vir + cur.sv; }
template<typename T>
                                                              bool connected(int u, int v) {
using ordered_set = tree<T, null_type, less<T>,
                                                               return find_root(++u) == find_root(++v); }
   rb_tree_tag, tree_order_statistics_node_update>;
                                                              void link(int x, int y) {
  find_by_order, order_of_key
                                                               chroot(++x); access(++y);
// hash tables: cc_hash_table/gp_hash_table
                                                               o[y].vir = o[y].vir + o[x].sub; up(o[x].pa = y);
2.2 Link-Cut Tree [2aaa19] - 0d97f7/f05d4f/642331
                                                              void cut(int x, int y) {
template <typename Val, typename SVal> class LCT {
                                                               split(++x, ++y); o[y].ch[0] = o[x].pa = 0; up(y); }
struct node {
 int pa, ch[2]; bool rev;
                                                             #undef cur
                                                             #undef lc
 Val v, prod, rprod; SVal sv, sub, vir;
 node() : pa{0}, ch{0, 0}, rev{false}, v{},
                                                             #undef rc
  prod{}, rprod{}, sv{}, sub{}, vir{} {}
                                                                   LiChao Segtree [8e1eaf]
                                                             // cmp(l, r, i) := is l better than r at i?
template <typename L, typename Cmp> class LiChao {
#define cur o[u]
#define lc cur.ch[0]
                                                              int n; vector<L> T; Cmp cmp;
void insert(int l, int r, int o, L ln) {
#define rc cur.ch[1]
vector<node> o;
bool is_root(int u) const {
                                                                // if (ln is empty line) return; // constant
  return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u; }
                                                               int m = (l + r) >> 1;
bool is_rch(int u) const {
                                                               bool atL = cmp(ln, T[o], l);
  return o[cur.pa].ch[1] == u && !is_root(u); }
                                                               if (cmp(ln, T[o], m)) atL ^= 1, swap(T[o], ln);
void down(int u) {
                                                               if (r - l == 1) return;
                                                               if (atL) insert(l, m, o << 1, ln);</pre>
 if (not cur.rev) return;
 for (int c : {lc, rc}) if (c) set_rev(c);
                                                               else insert(m, r, o << 1 | 1, ln);
 cur.rev = false;
                                                              L query(int x, int l, int r, int o) {
void up(int u) {
                                                               if (r - l == 1) return T[o];
                                                               int m = (l + r) >> 1;
 cur.prod = o[lc].prod * cur.v * o[rc].prod;
  cur.rprod = o[rc].rprod * cur.v * o[lc].rprod;
                                                               L s = (x < m ? query(x, l, m, o << 1)
                                                                 : query(x, m, r, o << 1 | 1));
 cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
                                                               return cmp(s, T[o], x) ? s : T[o];
void set_rev(int u) {
 swap(lc, rc), swap(cur.prod, cur.rprod);
cur.rev ^= 1;
                                                             public:
                                                              LiChao(int n_, L init, Cmp &&c) : n(n_{-}), T(n * 4, init)
}
                                                                 ), cmp(c) {}
 /* SPLIT_HASH_HERE */
                                                              void insert(L ln) { insert(0, n, 1, ln); }
void rotate(int u) {
                                                              L query(int x) { return query(x, 0, n, 1); }
 int f = cur.pa, g = o[f].pa, l = is_rch(u);
if (cur.ch[l ^ 1]) o[cur.ch[l ^ 1]].pa = f;
                                                             };
// struct Line { lld a, b; };
  if (not is_root(f)) o[g].ch[is_rch(f)] = u;
                                                             // LiChao lct(
 o[f].ch[l] = cur.ch[l ^ 1], cur.ch[l ^ 1] = f;
                                                                 int(xs.size()), Line{0, INF},
  cur.pa = g, o[f].pa = u; up(f);
                                                                 [&u](const Line &l, const Line &r, int i) {
                                                                  lld x = xs[i];
                                                                  return l.a * x + l.b < r.a * x + r.b;
void splay(int u) {
                                                                 Treap* [ae576c]
 vector<int> stk = {u};
 while (not is_root(stk.back()))
                                                              _gnu_cxx::sfmt19937 rnd(7122); // <ext/random>
   stk.push_back(o[stk.back()].pa);
                                                             namespace Treap {
 while (not stk.empty())
  down(stk.back()), stk.pop_back();
                                                             struct node {
                                                              int size, pri; node *lc, *rc, *pa;
node() : size(1), pri(rnd()), lc(0), rc(0), pa(0) {}
  for (int f = cur.pa; not is_root(u); f = cur.pa) {
   if (!is_root(f))
   rotate(is_rch(u) == is_rch(f) ? f : u);
                                                              void pull() {
  rotate(u);
                                                               size = 1; pa = 0;
                                                               if (lc) { size += lc->size; lc->pa = this;
                                                               if (rc) { size += rc->size; rc->pa = this; }
 up(u);
void access(int x) {
                                                             };
  for (int u = x, last = 0; u; u = cur.pa) {
                                                             int SZ(node *x) { return x ? x->size : 0; }
  splay(u);
                                                             node *merge(node *L, node *R) {
                                                              if (not L or not R) return L ? L : R;
   cur.vir = cur.vir + o[rc].sub - o[last].sub;
                                                              if (L->pri > R->pri)
   rc = last; up(last = u);
 }
                                                               return L->rc = merge(L->rc, R), L->pull(), L;
  splay(x);
                                                              else
                                                               return R->lc = merge(L, R->lc), R->pull(), R;
```

```
void splitBySize(node *o, int k, node *&L, node *&R) {
                                                                    Graph
 if (not o) L = R = 0;
                                                                    SCC [16c7d6]
                                                               3.1
 else if (int s = SZ(o->lc) + 1; s <= k)
                                                               class SCC { // test @ library checker
  L=o, splitBySize(o->rc, k-s, L->rc, R), L->pull();
                                                               protected:
 else
                                                                int n, dfc, nscc; vector<vector<int>> G;
  R=o, splitBySize(o->lc, k, L, R->lc), R->pull();
                                                                vector<int> vis, low, idx, stk;
void dfs(int i) {
} // SZ(L) == k
int getRank(node *o) { // 1-base
                                                                 vis[i] = low[i] = ++dfc; stk.push_back(i);
 int r = SZ(o->lc) + 1;
                                                                 for (int j : G[i])
  if (!vis[j])
 for (; o->pa; o = o->pa)
  if (o->pa->rc == o) r += SZ(o->pa->lc) + 1;
                                                                   dfs(j), low[i] = min(low[i], low[j]);
 return r;
                                                                  else if (vis[j] != -1)
                                                                   low[i] = min(low[i], vis[j]);
} // namespace Treap
                                                                 if (low[i] == vis[i])
                                                                  for (idx[i] = nscc++; vis[i] != -1;) {
2.5 Linear Basis* [138d5d]
                                                                   int x = stk.back(); stk.pop_back();
template <int BITS, typename S = int> struct Basis {
  static constexpr S MIN = numeric_limits<S>::min();
                                                                   idx[x] = idx[i]; vis[x] = -1;
 array<pair<llu, S>, BITS> b;
                                                                }
 Basis() { b.fill({0, MIN}); }
                                                               public:
 void add(llu x, S p) {
                                                                SCC(int n_{-}) : n(n_{-}), dfc(0), nscc(0), G(n),
  for (int i = BITS-1; i>=0; i--) if (x >> i & 1) {
                                                                 vis(n), low(n), idx(n) {}
   if (b[i].first == 0) return b[i]={x, p}, void();
                                                                void add_edge(int u, int v) { G[u].push_back(v); }
   if (b[i].second < p)</pre>
                                                                void solve() {
    swap(b[i].first, x), swap(b[i].second, p);
                                                                 for (int i = 0; i < n; i++) if (!vis[i]) dfs(i); }</pre>
   x ^= b[i].first;
                                                                int get_id(int x) { return idx[x]; }
  }
                                                                int count() { return nscc; }
                                                               }; // dag edges point from idx large to idx small
 optional<llu> query_kth(llu v, llu k) {
                                                               3.2 2-SAT [ca961f]
  vector<pair<llu, int>> o;
for (int i = 0; i < BITS; i++)</pre>
                                                               struct TwoSat : SCC {
                                                                void orr(int x, int y) {
   if (b[i].first) o.emplace_back(b[i].first, i);
                                                                 if ((x ^ y) == 1) return;
  if (k >= (1ULL << o.size())) return {};
                                                                 add_edge(x ^ 1, y); add_edge(y ^ 1, x);
  for (int i = int(o.size()) - 1; i >= 0; i--)
   if ((k >> i & 1) ^ (v >> o[i].second & 1))
                                                                vector<int> solve2sat() {
    v ^= o[i].first;
                                                                 solve(); vector<int> res(n);
  return v;
                                                                 for (int i = 0; i < n; i += 2)
  if (idx[i] == idx[i + 1]) return {};</pre>
 Basis filter(S l) {
                                                                 for (int i = 0; i < n; i++)</pre>
  Basis res = *this;
                                                                  res[i] = idx[i] < idx[i ^ 1];
  for (int i = 0; i < BITS; i++)</pre>
                                                                 return res;
   if (res.b[i].second < l) res.b[i] = {0, MIN};</pre>
                                                                }
  return res;
                                                               3.3
                                                                     BCC [6ac6db]
};
                                                               class BCC {
                                                                int n, ecnt, bcnt;
      Binary Search on Segtree [6c61c0]
                                                                vector<vector<pair<int, int>>> g;
// find_first = l \rightarrow minimal \times s.t. check([l, x))
                                                                vector<int> dfn, low, bcc, stk;
   find_{last} = r \rightarrow maximal x s.t. check([x, r))
                                                                vector<bool> ap, bridge;
int find_first(int l, auto &&check) {
                                                                void dfs(int u, int f)
 if (l >= n) return n + 1;
                                                                 dfn[u] = low[u] = dfn[f] + 1;
 l += sz; push(l); Monoid sum; // identity
                                                                 int ch = 0;
                                                                 for (auto [v, t] : g[u]) if (bcc[t] == -1) {
  while ((l & 1) == 0) l >>= 1;
                                                                  bcc[t] = 0; stk.push_back(t);
  if (auto s = sum + nd[l]; check(s)) {
                                                                  if (dfn[v]) {
   while (l < sz) {</pre>
                                                                   low[u] = min(low[u], dfn[v]);
    prop(l); l = (l << 1);
                                                                   continue;
    if (auto nxt = sum + nd[l]; not check(nxt))
     sum = nxt, l++;
                                                                  ++ch, dfs(v, u);
                                                                  low[u] = min(low[u], low[v]);
if (low[v] > dfn[u]) bridge[t] = true;
   return l + 1 - sz;
  } else sum = s, l++;
                                                                  if (low[v] < dfn[u]) continue;</pre>
 } while (lowbit(l) != l);
                                                                  ap[u] = true;
 return n + 1;
                                                                  while (not stk.empty()) {
                                                                   int o = stk.back(); stk.pop_back();
int find_last(int r, auto &&check) {
                                                                   bcc[o] = bcnt;
 if (r <= 0) return -1;
                                                                   if (o == t) break;
 r += sz; push(r - 1); Monoid sum; // identity
 do {
                                                                  bcnt += 1:
  while (r > 1 and (r & 1)) r >>= 1;
                                                                 ap[u] = ap[u] and (ch != 1 or u != f);
  if (auto s = nd[r] + sum; check(s)) {
   while (r < sz) {</pre>
                                                               public:
    prop(r); r = (r << 1) | 1;
                                                                BCC(int n_{-}) : n(n_{-}), ecnt(0), bcnt(0), g(n), dfn(n),
    if (auto nxt = nd[r] + sum; not check(nxt))
                                                                   low(n), stk(), ap(n) {}
     sum = nxt, r--;
                                                                void add_edge(int u, int v) {
                                                                 g[u].emplace_back(v, ecnt);
   return r - sz;
                                                                 g[v].emplace_back(u, ecnt++);
  } else sum = s;
                                                                }
 } while (lowbit(r) != r);
                                                                void solve() {
 return -1;
```

bridge.assign(ecnt, false); bcc.assign(ecnt, -1);

```
for (int i = 0; i < n; ++i) if (!dfn[i]) dfs(i, i);</pre>
                                                                 vector<int> solve(int S, int T) {
int bcc_id(int x) const { return bcc[x]; }
                                                                  g[S].insert(g[S].begin(), T); dfs(S);
                                                                  vector<int> nxt(n + 1, n), prv = nxt;
nxt[S] = T; prv[T] = S; sgn[S] = -1;
bool is_ap(int x) const { return ap[x]; }
bool is_bridge(int x) const { return bridge[x]; }
                                                                  for (int i : ord) if (i != S && i != T) {
                                                                   int p = pa[i], l = ord[low[i] - 1];
    Round Square Tree [cf6d74]
                                                                   if (sgn[l] > 0) // insert after
struct RST { // be careful about isolate point
                                                                    nxt[i] = nxt[prv[i] = p], nxt[p] = prv[nxt[p]] = i;
 int n; vector<vector<int>> T;
RST(auto &G) : n(int(G.size())), T(n) {
                                                                    prv[i] = prv[nxt[i] = p], prv[p] = nxt[prv[p]] = i;
 vector<int> stk, vis(n), low(n);
auto dfs = [&](auto self, int u, int d) -> void {
                                                                   sgn[p] = -sgn[l];
   low[u] = vis[u] = d; stk.push_back(u);
                                                                  vector<int> v;
   for (int v : G[u]) if (!vis[v]) {
                                                                  for (int x = S; x != n; x = nxt[x]) v.push_back(x);
    self(self, v, d + 1);
if (low[v] == vis[u]) {
                                                                  return v;
                                                                 } // S, T are unique source / unique sink
     int cnt = int(T.size()); T.emplace_back();
                                                                 void add_edge(int a, int b) {
     for (int x = -1; x != v; stk.pop_back())
                                                                  g[a].emplace_back(b); g[b].emplace_back(a); }
      T[cnt].push_back(x = stk.back());
                                                                      存在 ST 雙極定向 iff 連接 (S,T) 後整張圖點雙連通
DMST [f4317e]
     T[u].push_back(cnt); // T is rooted
    } else low[u] = min(low[u], low[v]);
                                                                using lld = int64_t;
   } else low[u] = min(low[u], vis[v]);
                                                                struct E { int s, t; lld w; }; // 0-base
 };
                                                                struct PQ {
 for (int u = 0; u < n; u++)
                                                                 struct P {
if (!vis[u]) dfs(dfs, u, 1);
} // T may be forest; after dfs, stk are the roots
                                                                  lld v; int i;
                                                                  bool operator>(const P &b) const { return v > b.v; }
}; // test @ 2020 Shanghai K
3.5 Edge TCC [5a2668]
                                                                 win_heap<P> pq; lld tag;
void push(P p) { p.v -= tag; pq.emplace(p); }
P top() { P p = pq.top(); p.v += tag; return p; }
vector<vector<int>> ETCC(auto &adj) {
const int n = static_cast<int>(adj.size());
vector<int> up(n), low(n), in, out, nx, id;
                                                                 void join(PQ &b) {
in = out = nx = id = vector < int > (n, -1);
                                                                  if (pq.size() < b.pq.size())</pre>
int dfc = 0, cnt = 0; Dsu dsu(n);
auto merge = [&](int u, int v) {
                                                                   swap(pq, b.pq), swap(tag, b.tag);
                                                                  while (!b.pq.empty()) push(b.top()), b.pq.pop();
 dsu.join(u, v); up[u] += up[v]; };
auto dfs = [&](auto self, int u, int p) -> void {
                                                                };
  in[u] = low[u] = dfc++
                                                                vector<int> dmst(const vector<E> &e, int n, int root) {
  for (int v : adj[u]) if (v != u) {
                                                                 vector<PQ> h(n * 2);
for (int i = 0; i < int(e.size()); ++i)</pre>
  if (v == p) { p = -1; continue; }
if (in[v] == -1) {
                                                                  h[e[i].t].push({e[i].w, i});
    self(self, v, u);
if (nx[v] == -1 && up[v] <= 1) {</pre>
                                                                 vector<int> a(n * 2); iota(all(a), 0);
vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
     up[u] += up[v]; low[u] = min(low[u], low[v]);
                                                                 auto o = [\&](auto Y, int X) \rightarrow int {
     continue:
                                                                  return x==a[x] ? x : a[x] = Y(Y, a[x]); };
                                                                 auto S = [&](int i) { return o(o, e[i].s); };
    if (up[v] == 0) v = nx[v];
                                                                 int pc = v[root] = n;
    if (low[u] > low[v])
                                                                 for (int i = 0; i < n; ++i) if (v[i] == -1)
     low[u] = low[v], swap(nx[u], v);
                                                                  for (int p = i; v[p]<0 || v[p]==i; p = S(r[p])) {</pre>
  for (; v != -1; v = nx[v]) merge(u, v);
} else if (in[v] < in[u]) {</pre>
                                                                   if (v[p] == i)
                                                                     for (int q = pc++; p != q; p = S(r[p])) {
    low[u] = min(low[u], in[v]); up[u]++;
                                                                     h[p].tag -= h[p].top().v; h[q].join(h[p]);
   } else {
                                                                     pa[p] = a[p] = q;
    for (int &x = nx[u]; x != -1 &&
      in[x] \le in[v] \&\& in[v] \le out[x]; x = nx[x])
                                                                   while (S(h[p].top().i) == p) h[p].pq.pop();
     merge(u, x);
                                                                   v[p] = i; r[p] = h[p].top().i;
    up[u]--;
  }
                                                                 vector<int> ans;
                                                                 for (int i = pc - 1; i >= 0; i--) if (v[i] != n) {
  out[u] = dfc;
                                                                  for (int f = e[r[i]].t; f!=-1 && v[f]!=n; f = pa[f])
};
                                                                   v[f] = n:
for (int i = 0; i < n; i++)</pre>
                                                                  ans.push_back(r[i]);
  if (in[i] == -1) dfs(dfs, i, -1);
 for (int i = 0; i < n; i++)
                                                                 return ans; // default minimize, returns edgeid array
 if (dsu.anc(i) == i) id[i] = cnt++;
 vector<vector<int>> comps(cnt);
                                                                3.8 Dominator Tree [ea5b7c]
for (int i = 0; i < n; i++)
                                                                struct Dominator {
  comps[id[dsu.anc(i)]].push_back(i);
                                                                 vector<vector<int>> g, r, rdom; int tk;
vector<int>> dfn, rev, fa, sdom, dom, val, rp;
return comps;
} // test @ yosupo judge
                                                                 Dominator(int n) : g(n), r(n), rdom(n), tk(0) {
3.6 Bipolar Orientation [b50cd3]
                                                                  dfn = rev = fa = sdom = dom =
                                                                   val = rp = vector<int>(n, -1); }
struct BipolarOrientation {
int n; vector<vector<int>> g;
                                                                 void add_edge(int x, int y) { g[x].push_back(y); }
vector<int> vis, low, pa, sgn, ord;
                                                                 void dfs(int x) {
                                                                  rev[dfn[x] = tk] = x;
BipolarOrientation(int n_) : n(n_),
g(n), vis(n), low(n), pa(n, -1), sgn(n) {} void dfs(int i) {
                                                                  fa[tk] = sdom[tk] = val[tk] = tk; tk++;
                                                                  for (int u : g[x]) {
                                                                   if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
 ord.push_back(i); low[i] = vis[i] = int(ord.size());
  for (int j : g[i])
                                                                   r[dfn[u]].push_back(dfn[x]);
   if (!vis[j])
                                                                  }
    pa[j] = i, dfs(j), low[i] = min(low[i], low[j]);
   else low[i] = min(low[i], vis[j]);
                                                                 void merge(int x, int y) { fa[x] = y; }
```

```
int find(int x, int c = 0) {
                                                                3.10 Centroid Decomp.* [670cdd]
  if (fa[x] == x) return c ? -1 : x;
  if (int p = find(fa[x], 1); p != -1) {
                                                                 class Centroid {
   if (sdom[val[x]] > sdom[val[fa[x]]])
                                                                  vector<vector<pair<int, int>>> g; // g[u] = {(v, w)}
    val[x] = val[fa[x]];
                                                                  vector<int> pa, dep, vis, sz, mx;
   fa[x] = p;
                                                                  vector<vector<int64_t>> Dist;
   return c ? p : val[x];
                                                                  vector<int64_t> Sub, Sub2;
                                                                  vector<int> Cnt, Cnt2;
  } else return c ? fa[x] : val[x];
                                                                  void DfsSz(vector<int> &tmp, int x) {
vector<int> build(int s, int n) {
// return the father of each node in dominator tree
                                                                   vis[x] = true, sz[x] = 1, mx[x] = 0;
for (auto [u, w] : g[x]) if (not vis[u]) {
 dfs(s); // p[i] = -2 if i is unreachable from s
                                                                    DfsSz(tmp, u); sz[x] += sz[u];
  for (int i = tk - 1; i >= 0; --i) {
                                                                    mx[x] = max(mx[x], sz[u]);
   for (int u : r[i])
    sdom[i] = min(sdom[i], sdom[find(u)]);
                                                                   tmp.push_back(x);
   if (i) rdom[sdom[i]].push_back(i);
   for (int u : rdom[i]) {
                                                                  void DfsDist(int x, int64_t D = 0) {
    int p = find(u);
                                                                   Dist[x].push_back(D); vis[x] = true;
    dom[u] = (sdom[p] == i ? i : p);
                                                                   for (auto [u, w] : g[x])
                                                                    if (not vis[u]) DfsDist(u, D + w);
   if (i) merge(i, rp[i]);
  }
                                                                  void DfsCen(int x, int D, int p) {
                                                                   vector<int> tmp; DfsSz(tmp, x);
int M = int(tmp.size()), C = -1;
 vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)</pre>
   if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
                                                                   for (int u : tmp)
 for (int i = 1; i < tk; ++i)
                                                                    if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
   p[rev[i]] = rev[dom[i]];
                                                                   for (int u : tmp) vis[u] = false;
  return p;
                                                                   DfsDist(C);
} // test @ yosupo judge
                                                                   for (int u : tmp) vis[u] = false;
                                                                   pa[C] = p, vis[C] = true, dep[C] = D;
      Edge Coloring [029763]
                                                                   for (auto [u, w] : g[C])
                                                                    if (not vis[u]) DfsCen(u, D + 1, C);
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
                                                                 public:
                                                                  Centroid(int N) : g(N), pa(N), dep(N),
vis(N), sz(N), mx(N), Dist(N),
Sub(N), Sub2(N), Cnt(N), Cnt2(N) {}
for (int i = 0; i <= N; i++)</pre>
  for (int j = 0; j <= N; j++)</pre>
    C[i][j] = G[i][j] = 0;
                                                                  void AddEdge(int u, int v, int w) {
                                                                   g[u].emplace_back(v, w);
void solve(vector<pair<int, int>> &E, int N) {
int X[kN] = {}, a;
                                                                   g[v].emplace_back(u, w);
auto update = [&](int u) {
 for (X[u] = 1; C[u][X[u]]; X[u]++);
                                                                  void Build() { DfsCen(0, 0, -1); }
                                                                  void Mark(int v) {
auto color = [&](int u, int v, int c) {
                                                                   int x = v, z = -1;
                                                                   for (int i = dep[v]; i >= 0; --i) {
 int p = G[u][v];
 G[u][v] = G[v][u] = c;
                                                                    Sub[x] += Dist[v][i], Cnt[x]++;
                                                                    if (z != -1)
 C[u][c] = v, C[v][c] = u;
                                                                     Sub2[z] += Dist[v][i], Cnt2[z]++;
 C[u][p] = C[v][p] = 0;
  if (p) X[u] = X[v] = p;
                                                                    x = pa[z = x];
 else update(u), update(v);
                                                                   }
  return p;
};
                                                                  int64_t Query(int v) {
auto flip = [&](int u, int c1, int c2) {
                                                                   int64_t res = 0;
                                                                   int x = v, z = -1;
for (int i = dep[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Cnt[x] * Dist[v][i];
 int p = C[u][c1];
 swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
 if (!C[u][c1]) X[u] = c1;
                                                                    if (z != -1)
 if (!C[u][c2]) X[u] = c2;
                                                                     res -= Sub2[z] + 1LL * Cnt2[z] * Dist[v][i];
  return p;
                                                                    x = pa[z = x];
                                                                   }
};
for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {</pre>
                                                                   return res;
                                                                 3: // pa, dep are centroid tree attributes
3.11 Lowbit Decomp. [2d7032]
 auto [u, v] = E[t];
  int v0 = v, c = X[u], c0 = c, d;
  vector<pair<int, int>> L; int vst[kN] = {};
                                                                 class LBD {
 while (!G[u][v0]) {
                                                                  int n, timer, chains;
   L.emplace_back(v, d = X[v]);
if (!C[v][c]) for (a=L.size()-1;a>=0;a--)
                                                                  vector<vector<int>> G;
                                                                  vector<int> tl, tr, chain, top, dep, pa;
// chains : number of chain
     c = color(u, L[a].first, c);
   else if (!C[u][d]) for (a=L.size()-1;a>=0;a--)
                                                                  // tl, tr[u] : subtree interval in the seq. of u
     color(u, L[a].first, L[a].second);
                                                                  // top[i] : top of the chain of vertex i
                                                                  // chian[u] : chain id of the chain u is on
   else if (vst[d]) break
   else vst[d] = 1, v = C[u][d];
                                                                  void predfs(int u, int f) {
                                                                   dep[u] = dep[pa[u] = f] + 1;
  if (!G[u][v0]) {
                                                                   for (int v : G[u]) if (v != f) {
   for (; v; v = flip(v, c, d), swap(c, d));
                                                                    predfs(v, u);
   if (C[u][c0]) { a = int(L.size()) - 1;
                                                                    if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
    while (--a >= 0 && L[a].second != c);
                                                                     chain[u] = chain[v];
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
   } else t--;
                                                                   if (chain[u] == 0) chain[u] = ++chains;
                                                                  void dfschain(int u, int f, int t) {
```

```
tl[u] = timer++; top[u] = t;
  for (int v : G[u])
  if (v != f and chain[v] == chain[u])
   dfschain(v, u, t);
  for (int v : G[u])
  if (v != f and chain[v] != chain[u])
   dfschain(v, u, v);
 tr[u] = timer;
                                                            private:
public:
LBD(auto &&G_) : n((int)size(G_)),
 timer(0), chains(0), G(G_), tl(n), tr(n),
  chain(n), top(n + 1, -1), dep(n), pa(n)
  { predfs(0, 0); dfschain(0, 0, 0);
PII get_subtree(int u) { return {tl[u], tr[u]}; }
vector<PII> get_path(int u, int v) {
 vector<PII> res;
 while (top[u] != top[v]) {
  if (dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
  int s = top[u];
  res.emplace_back(tl[s], tl[u] + 1);
  u = pa[s];
 if (dep[u] < dep[v]) swap(u, v);</pre>
 res.emplace_back(tl[v], tl[u] + 1);
 return res;
                                                                return;
}; // 記得在資結上對點的修改要改成對其 dfs 序的修改
3.12 Virtual Tree* [44f764]
vector<pair<int, int>> build(vector<int> vs, int r) {
vector<pair<int, int>> res;
sort(vs.begin(), vs.end(), [](int i, int j) {
 return dfn[i] < dfn[j]; });</pre>
vector<int> s = {r};
for (int v : vs) if (v != r) {
 if (int o = lca(v, s.back()); o != s.back()) {
  while (s.size() >= 2) {
                                                              }
                                                              }
   if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
                                                            public:
    res.emplace_back(s[s.size() - 2], s.back());
   s.pop_back();
  if (s.back() != o)
   res.emplace_back(o, s.back()), s.back() = o;
 s.push_back(v);
                                                              int solve() {
for (size_t i = 1; i < s.size(); ++i)</pre>
 res.emplace_back(s[i - 1], s[i]);
return res; // (x, y): x->y
} // 記得建虛樹會多出 `vs` 以外的點
3.13 Tree Hashing [d6a9f9]
vector<int> g[maxn]; llu h[maxn];
llu F(llu z) { // xorshift64star from iwiwi
z \stackrel{\wedge}{=} z >> 12; z \stackrel{\wedge}{=} z << 25; z \stackrel{\wedge}{=} z >> 27;
return z * 2685821657736338717LL;
                                                              }
llu hsah(int u, int f) {
llu r = 127; // bigger?
for (int v : g[u]) if (v != f) r += hsah(v, u);
return h[u] = F(r);
} // test @ UOJ 763 & yosupo library checker
3.14 Mo's Algo on Tree
dfs u:
push u
                                                              n = _n;
iterate subtree
push u
Let P = LCA(u, v) with St(u) \le St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
                                                               if (i < 4) {
3.15 Count Cycles [c7e8f2]
// ord = sort by deg decreasing, rk[ord[i]] = i
// D[i] = edge point from rk small to rk big
for (int x : ord) { // c3
for (int y : D[x]) vis[y] = 1;
for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
for (int y : D[x]) vis[y] = 0;
for (int x : ord) { // c4
for (int y : D[x]) for (int z : adj[y])
 if (rk[z] > rk[x]) c4 += vis[z]++;
for (int y : D[x]) for (int z : adj[y])
```

```
if (rk[z] > rk[x]) --vis[z];
// both are O(M*sqrt(M)), test @ 2022 CCPC guangzhou
3.16 Maximal Clique [2da556]
#define iter(u, B) for (size_t u = B._Find_first(); \
  u < n; u = B._Find_next(u))</pre>
// contain a self loop u to u, than u won't in clique
template <size_t maxn> class MaxClique {
 using bits = bitset<maxn>;
 bits popped, G[maxn], ans;
 size_t deg[maxn], deo[maxn], n;
 void sort_by_degree() {
  popped.reset();
  for (size_t i = 0; i < n; ++i) deg[i] = G[i].count();</pre>
  for (size_t i = 0; i < n; ++i) {</pre>
   size_t mi = maxn, id = 0;
   for (size_t j = 0; j < n; ++j)</pre>
    if (!popped[j] and deg[j] < mi) mi = deg[id = j];</pre>
   popped[deo[i] = id] = 1;
   iter(u, G[i]) --deg[u];
 void BK(bits R, bits P, bits X) {
  if (R.count() + P.count() <= ans.count()) return;</pre>
  if (not P.count() and not X.count()) {
   if (R.count() > ans.count()) ans = R;
  /* greedily chosse max degree as pivot
  bits cur = P \mid X; size_t pv = 0, sz = 0; iter(u, cur) if (deg[u] > sz) sz = deg[pv = u];
  cur = P \& \sim G[pv] \& \sim R; */ // or simply choose first
  bits cur = P & (~G[(P | X)._Find_first()]) & ~R;
  iter(u, cur) {
   R[u] = 1; BK(R, P \& G[u], X \& G[u]);
   R[u] = P[u] = 0, X[u] = 1;
 void init(size_t n_) {
  n = n_; ans.reset();
  for (size_t i = 0; i < n; ++i) G[i].reset();</pre>
 void add_edges(int u, bits S) { G[u] = S; }
 void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
  sort_by_degree(); // or simply iota( deo...
  for (size_t i = 0; i < n; ++i) deg[i] = G[i].count();</pre>
  bits pob, nob = 0; pob.set();
  for (size_t i = n; i < maxn; ++i) pob[i] = 0;</pre>
  for (size_t i = 0; i < n; ++i) {</pre>
   size_t v = deo[i]; bits tmp; tmp[v] = 1;
   BK(tmp, pob \& G[v], nob \& G[v]);
   pob[v] = 0, nob[v] = 1;
  return static cast<int>(ans.count()):
3.17 Maximum Clique [aee5d8]
constexpr size_t kN = 150; using bits = bitset<kN>;
struct MaxClique {
 bits G[kN], cs[kN];
 int ans, sol[kN], q, cur[kN], d[kN], n;
 void init(int _n) {
  for (int i = 0; i < n; ++i) G[i].reset();</pre>
 void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
 void pre_dfs(vector<int> &v, int i, bits mask) {
   for (int x : v) d[x] = (int)(G[x] \& mask).count();
   sort(all(v), [&](int x, int y) {
    return d[x] > d[y]; });
  vector<int> c(v.size());
  cs[1].reset(), cs[2].reset();
  int l = max(ans - q + 1, 1), r = 2, tp = 0, k;
  for (int p : v) {
   for (k = 1; (cs[k] & G[p]).any(); ++k);
   if (k >= r) cs[++r].reset();
   cs[k][p] = 1;
   if (k < l) v[tp++] = p;
```

```
for (k = 1; k < r; ++k)
   for (auto p = cs[k]._Find_first();
     p < kN; p = cs[k]._Find_next(p))</pre>
    v[tp] = (int)p, c[tp] = k, ++tp;
  dfs(v, c, i + 1, mask);
 void dfs(vector<int> &v, vector<int> &c,
   int i, bits mask) {
  while (!v.empty()) {
   int p = v.back(); v.pop_back(); mask[p] = 0;
   if (q + c.back() <= ans) return;</pre>
   cur[q++] = p;
   vector<int> nr;
   for (int x : v) if (G[p][x]) nr.push_back(x);
   if (!nr.empty()) pre_dfs(nr, i, mask & G[p]);
   else if (q > ans) ans = q, copy_n(cur, q, sol);
   c.pop_back(); --q;
 }
 int solve() {
  vector<int> v(n); iota(all(v), 0);
  ans = q = 0; pre_dfs(v, 0, bits(string(n, '1')));
  return ans; // sol[0 ~ ans-1]
  cliq; // test @ yosupo 1909
18 Min Mean Cycle [e23bc0]
3.18 <sup>1</sup>
// WARNING: TYPE matters
struct Edge { int s, t; llf c; };
llf solve(vector<Edge> &e, int n) {
 // O(VE), returns inf if no cycle, mmc otherwise
 vector<VI> prv(n + 1, VI(n)), prve = prv;
 vector<vector<llf>> d(n + 1, vector<llf>(n, inf));
 d[0] = vector<llf>(n, 0);
 for (int i = 0; i < n; i++) {</pre>
  for (int j = 0; j < (int)e.size(); j++) {
  auto [s, t, c] = e[j];</pre>
   if (d[i][s] < inf && d[i + 1][t] > d[i][s] + c) {
    d[i + 1][t] = d[i][s] + c;
    prv[i + 1][t] = s; prve[i + 1][t] = j;
  }
 llf mmc = inf; int st = -1;
 for (int i = 0; i < n; i++) {</pre>
  llf avg = -inf;
  for (int k = 0; k < n; k++) {
  if (d[n][i] < inf - eps)</pre>
    avg = max(avg, (d[n][i] - d[k][i]) / (n - k));
   else avg = inf;
  if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
 if (st == -1) return inf;
 vector<int> vst(n), eid, cycle, rho;
for (int i = n; !vst[st]; st = prv[i--][st]) {
  vst[st]++; eid.emplace_back(prve[i][st]);
  rho.emplace_back(st);
 while (vst[st] != 2) {
  int v = rho.back(); rho.pop_back();
  cycle.emplace_back(v); vst[v]++;
 reverse(all(eid)); eid.resize(cycle.size());
 return mmc;
3.19
     Eulerian Trail [8a70bf]
// g[i] = list of (edge.to, edge.id)
auto euler(int N, int M, int S, const auto &g) {
 vector<int> iter(N), vis(M), vv, ee;
 auto dfs = [&](auto self, int i) -> void {
  while (iter[i] < ssize(g[i])) {</pre>
   auto [j, eid] = g[i][iter[i]++];
   if (vis[eid]) continue;
   vis[eid] = true; self(self, j);
   vv.push_back(j); ee.push_back(eid);
  }
 dfs(dfs, S); vv.push_back(S);
 reverse(all(vv)); reverse(all(ee));
 return pair{vv, ee};
\} // 需要保證傳入的 g, S degree 符合條件;小心孤點奇點
```

```
4 Flow & Matching
4.1 HopcroftKarp [930040]
struct HK {
```

```
vector<int> l, r, a, p; int ans;
HK(int n, int m, auto &g) : l(n,-1), r(m,-1), ans(0) {
  for (bool match = true; match;) {
   match = false; a.assign(n, -1); p.assign(n, -1);
   queue<int> q;
   for (int i = 0; i < n; i++)</pre>
    if (l[i] == -1) q.push(a[i] = p[i] = i);
   // bitset<maxn> nvis, t; nvis.set();
   while (!q.empty()) {
    int z, x = q.front(); q.pop();
    if (l[a[x]] != -1) continue;
    for (int y : g[x]) { // or iterate t = g[x]&nvis
      // nvis.reset(y);
     if (r[y] == -1) {
      for (z = y; z != -1; )
r[z] = x, swap(l[x], z), x = p[x];
      match = true; ++ans; break;
     } else if (p[r[y]] == -1)
      q.push(z = r[y]), p[z] = x, a[z] = a[x];
  }
4.2
     Kuhn Munkres* [74bf6d]
struct KM { // maximize, test @ UOJ 80
 int n, l, r; lld ans; // fl and fr are the match
 vector<lid> hl, hr; vector<int> fl, fr, pre, q;
void bfs(const auto &w, int s) {
  vector<int> vl(n), vr(n); vector<lld> slk(n, INF);
  l = r = 0; vr[q[r++] = s] = true;
  auto check = [\&](int x) -> bool {
   if (vl[x] || slk[x] > 0) return true;
   vl[x] = true; slk[x] = INF;
   if (fl[x] != -1) return (vr[q[r++] = fl[x]] = true);
   while (x != -1) swap(x, fr[fl[x] = pre[x]]);
   return false;
  while (true) {
   while (l < r)
    for (int x = 0, y = q[l++]; x < n; ++x) if (!vl[x])
     if (chmin(slk[x], hl[x] + hr[y] - w[x][y]))
      if (pre[x] = y, !check(x)) return;
   lld d = ranges::min(slk);
   for (int x = 0; x < n; ++x)
    vl[x] ? hl[x] += d : slk[x] -= d;
   for (int x = 0; x < n; ++x) if (vr[x]) hr[x] -= d;
for (int x = 0; x < n; ++x) if (!check(x)) return;</pre>
```

// $K=1\ldots N$, initialize hl[i]=INF and bfs from all // unmatched right part point (fr[i]==-1)4.3 Flow Models

Maximum/Minimum flow with lower bound / Circulation problem

 $KM(int n_{, const auto \&w) : n(n_{, ans(0), ans(0),$

hl(n), hr(n), fl(n, -1), fr(fl), pre(n), q(n) { for (int i = 0; i < n; ++i) hl[i]=ranges::max(w[i]);
for (int i = 0; i < n; ++i) bfs(w, i);</pre>

for (int i = 0; i < n; ++i) ans += w[i][fl[i]];</pre>

// To obtain the max match of exactly K edges for

1. Construct super source ${\cal S}$ and sink ${\cal T}$.

}; // find maximum perfect matching

- 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
- For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
- 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect v o T with capacity -in(v). – To maximize, connect t o s with capacity ∞ (skip this in circu
 - lation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer. Also, f is a mincost valid flow.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.
- 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph. ullet Construct minimum vertex cover from maximum matching M on bipartite
- graph(X, Y)1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise. 2. DFS from unmatched vertices in X.

Cap max_flow(int st_, int ed_) {

st = st_, ed = ed_; Cap ret = 0;

```
x \in X is chosen iff x is unvisited; y \in Y is chosen iff y is visited.
                                                                                                              for (int i = 63; i >= 0; --i)

    Minimum cost cyclic flow

                                                                                                               while (BFS(i)) ret += DFS(st);
  1. Consruct super source {\cal S} and sink {\cal T}
                                                                                                              return ret;
  2. For each edge (x,y,c), connect x\to y with (cost,cap)=(c,1) if c>0, otherwise connect y\to x with (cost,cap)=(-c,1)
                                                                                                          }; // test @ luogu P3376
4.5 HI DD P07
  3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1,
      decrease d(\bar{x}) by 1
  4. For each vertex v with d(v) > 0, connect S \rightarrow v with (cost, cap) =
                                                                                                          template <typename T> struct HLPP {
  struct Edge { int to, rev; T flow, cap; };
       (0, d(v))
  5. For each vertex v with d(v) < 0, connect v \rightarrow T with (cost, cap) =
                                                                                                            int n, mx; vector<vector<Edge>> adj; vector<T> excess;
       (0, -d(v))
                                                                                                            vector<int> d, cnt, active; vector<vector<int>> B;
void add_edge(int u, int v, int f) {
  6. Flow from S to T, the answer is the cost of the flow C+K
· Maximum density induced subgraph
   1. Binary search on answer, suppose we're checking answer {\cal T}
                                                                                                              Edge a{v, (int)size(adj[v]), 0, f};
      Construct a max flow model, let K be the sum of all weights
                                                                                                              Edge b{u, (int)size(adj[u]), 0, 0};
  3. Connect source s \to v, v \in G with capacity K
4. For each edge (u,v,w) in G, connect u \to v and v \to u with capacity w
5. For v \in G, connect it with sink v \to t with capacity K + 2T - t
                                                                                                              adj[u].push_back(a), adj[v].push_back(b);
                                                                                                            void enqueue(int v) {
       \left(\sum_{e \in E(v)} w(e)\right) - 2w(v)
                                                                                                              if (!active[v] && excess[v] > 0 && d[v] < n) {</pre>
     \stackrel{.}{T} is a valid answer if the maximum flow f < K|V|
                                                                                                               mx = max(mx, d[v]);

    Minimum weight edge cover

                                                                                                               B[d[v]].push_back(v); active[v] = 1;
  1. For each v \in V create a copy v', and connect u' \to v' with weight
                                                                                                             }
  2. Connect v 	o v' with weight 2\mu(v), where \mu(v) is the cost of the cheap-
      est edge incident to v.
                                                                                                            void push(int v, Edge &e) {
   3. Find the minimum weight perfect matching on G^\prime
                                                                                                             T df = min(excess[v], e.cap - e.flow);
if (df <= 0 || d[v] != d[e.to] + 1) return;
 Project selection cheat sheet: S,T 分別代表 0,1 側,最小化總花費。
    i 為 O 時花費 c
                                              (i, T, c)
                                                                                                              e.flow += df, adj[e.to][e.rev].flow -= df;
    i 為 1 時花費 c
                                              (S, i, c)
    i \in I 有任何一個為 0 時花費 c i \in I 有任何一個為 1 時花費 c
                                                                                                              excess[e.to] += df, excess[v] -= df;
                                              (i, w, \infty), (w, T, c)
                                              (S, w, c), (w, i, \infty)
直接得到 c ; (S, i, c)
                                                                                                              enqueue(e.to);
    i 為 0 時得到 c
     i 為 1 時得到 c
                                              直接得到 c; (i, T, c)
                                                                                                            void gap(int k) {
    i 為 0,j 為 1 時花費 c
                                              (i, j, c)
                                             (i,j,c),(j,i,c)
直接得到 c ; (S,w,c),(w,i,\infty),(w,j,\infty)
直接得到 c ; (i,w,\infty),(j,w,\infty),(w,T,c)
                                                                                                              for (int v = 0; v < n; v++) if (d[v] >= k)
    i,j 不同時花費 c
    i,j 同時是 0 時得到 c
                                                                                                               cnt[d[v]]--, d[v] = n, cnt[d[v]]++;
     i, j 同時是 1 時得到 c
· Submodular functions minimization
                                                                                                            void relabel(int v) {
  - For a function f: 2^V \to \mathbb{R}, f is a submodular function iff * \forall S, T \subseteq V, f(S) + f(T) \geq f(S \cup T) + f(S \cap T), or * \forall X \subseteq Y \subseteq V, x \notin Y, f(X \cup \{x\}) - f(X) \geq f(Y \cup \{x\}) - f(Y). - To minimize \sum_i \theta_i(x_i) + \sum_{i < j} \phi_{ij}(x_i, x_j) + \sum_{i < j < k} \psi_{ijk}(x_i, x_j, x_k) - If \theta_i(1) \geq \theta_i(0), add edge \{S, i, \theta_i(1) - \theta_i(0)\} and \theta_i(0) to answer; otherwise f(X, X) = f(X).
                                                                                                             cnt[d[v]]--; d[v] = n;
                                                                                                              for (auto e : adj[v])
                                                                                                               if (e.cap > e.flow) d[v] = min(d[v], d[e.to] + 1);
                                                                                                              cnt[d[v]]++; enqueue(v);
  - Add edges (i,j,\phi_{ij}(0,1)+\phi_{ij}(1,0)-\phi_{ij}(0,0)-\phi_{ij}(1,1)).

- Denote x_{ijk} as helper nodes. Let P=\psi_{ijk}(0,0,0)+\psi_{ijk}(0,1,1)+\psi_{ijk}(1,0,1)+\psi_{ijk}(1,1,0)-\psi_{ijk}(0,0,1)-\psi_{ijk}(0,1,0)-\psi_{ijk}(1,1,1). Add -P to answer. If P\geq 0, add edges (i,x_{ijk},P), (j,x_{ijk},P), (k,x_{ijk},P), (x_{ijk},T,P); otherwise (x_{ijk},i,-P), (x_{ijk},j,-P), (x_{ijk},-P), (x_{i
     wise, (i, T, \theta_i(0) - \theta_i(1)) and \theta_i(1).
                                                                                                            void discharge(int v) {
                                                                                                              for (auto &e : adj[v])
                                                                                                                if (excess[v] > 0) push(v, e);
                                                                                                               else break:
                                                                                                              if (excess[v] <= 0) return;</pre>
   The minimum cut of this graph will be the the minimum value of the
                                                                                                              if (cnt[d[v]] == 1) gap(d[v]);
     function above
                                                                                                              else relabel(v);
4.4 Dinic [32c53e]
template <typename Cap = int64_t> class Dinic {
                                                                                                            T max_flow(int s, int t) {
                                                                                                              for (auto &e : adj[s]) excess[s] += e.cap;
private:
 struct E { int to, rev; Cap cap; }; int n, st, ed;
                                                                                                              cnt[0] = n; enqueue(s); active[t] = 1;
 vector<vector<E>> G; vector<size_t> lv, idx;
                                                                                                              for (mx = 0; mx >= 0;)
                                                                                                               if (!B[mx].empty()) {
 bool BFS(int k) {
   lv.assign(n, 0); idx.assign(n, 0);
                                                                                                                 int v = B[mx].back(); B[mx].pop_back();
   queue<int> bfs; bfs.push(st); lv[st] = 1;
                                                                                                                 active[v] = 0; discharge(v);
   while (not bfs.empty() and not lv[ed]) {
                                                                                                                } else --mx;
    int u = bfs.front(); bfs.pop();
for (auto e: G[u]) if (e.cap >> k and !lv[e.to])
                                                                                                              return excess[t];
      bfs.push(e.to), lv[e.to] = lv[u] + 1;
                                                                                                            HLPP(int _n) : n(_n), adj(n), excess(n),
                                                                                                             d(n), cnt(n + 1), active(n), B(n) {}
   }
   return lv[ed];
                                                                                                          4.6
                                                                                                                     Global Min-Cut [ae7013]
 Cap DFS(int u, Cap f = numeric_limits<Cap>::max()) {
                                                                                                          void add_edge(auto &w, int u, int v, int c) {
   w[u][v] += c; w[v][u] += c; }
   if (u == ed) return f;
   Cap ret = 0;
                                                                                                          auto phase(const auto &w, int n, vector<int> id) {
   for (auto &i = idx[u]; i < G[u].size(); ++i) {</pre>
                                                                                                            vector<lld> g(n); int s = -1, t = -1;
     auto &[to, rev, cap] = G[u][i];
                                                                                                            while (!id.empty()) {
     if (cap <= 0 or lv[to] != lv[u] + 1) continue;</pre>
                                                                                                              int c = -1;
    Cap nf = DFS(to, min(f, cap));
ret += nf; cap -= nf; f -= nf;
G[to][rev].cap += nf;
                                                                                                              for (int i : id) if (c == -1 || g[i] > g[c]) c = i;
                                                                                                              s = t; t = c;
                                                                                                              id.erase(ranges::find(id, c));
     if (f == 0) return ret;
                                                                                                              for (int i : id) g[i] += w[c][i];
   if (ret == 0) lv[u] = 0;
                                                                                                            return tuple{s, t, g[t]};
   return ret;
                                                                                                          lld mincut(auto w, int n) {
public:
                                                                                                            lld cut = numeric_limits<lld>::max();
                                                                                                            vector<int> id(n); iota(all(id), 0);
for (int i = 0; i < n - 1; ++i) {</pre>
 void init(int n_) { G.assign(n = n_, vector<E>()); }
 void add_edge(int u, int v, Cap c) {
G[u].push_back({v, int(G[v].size()), c});
                                                                                                              auto [s, t, gt] = phase(w, n, id);
   G[v].push_back({u, int(G[u].size())-1, 0});
                                                                                                              id.erase(ranges::find(id, t));
```

cut = min(cut, gt);

for (int j = 0; j < n; ++j)</pre>

w[s][j] += w[t][j], w[j][s] += w[j][t];

```
return cut;
                                                              if (d[T] == INF_C) return nullopt;
                                                              for (size_t i = 0; i < d.size(); ++i) h[i] += d[i];</pre>
for (int i = T; i != S; i = f[i].first) {
      GomoryHu Tree [245ce3]
                                                               auto &eg = g[f[i].first][f[i].second];
auto GomoryHu(int n, const auto &flow) {
                                                               eg.f -= up[T]; g[eg.to][eg.r].f += up[T];
vector<tuple<int, int, int>> rt; vector<int> g(n);
for (int i = 1; i < n; ++i) {
                                                              return pair{up[T], h[T]};
 int t = g[i]; auto f = flow;
  rt.emplace_back(f.max_flow(i, t), i, t);
                                                            public:
  f.walk(i); // bfs from i use edges with .cap > 0
                                                             MCMF(int n) : g(n), f(n), up(n), d(n, INF_C) {}
  for (int j = i + 1; j < n; ++j)</pre>
                                                             void add_edge(int s, int t, F c, C w) {
   if (g[j]==t && f.connect(j)) g[j] = i;
                                                              g[s].emplace_back(t, int(g[t].size()), c, w);
                                                              g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
return rt;
} // for our dinic:
                                                             pair<F, C> solve(int a, int b) {
// void walk(int) { BFS(0); }
                                                              h.assign(g.size(), 0);
// bool connect(int i) { return lv[i]; }
                                                              F c = 0; C w = 0;
4.8 MCMF [0df510]
                                                              while (auto r = step(a, b)) {
template <typename F, typename C> class MCMF {
    static constexpr F INF_F = numeric_limits<F>::max();
                                                                c += r->first, w += r->first * r->second;
                                                               fill(d.begin(), d.end(), INF_C);
static constexpr C INF_C = numeric_limits<C>::max();
struct E { int to, r; F f; C c; };
                                                              return {c, w};
vector<vector<E>> g; vector<pair<int, int>> f;
                                                             7
vector<int> inq; vector<F> up; vector<C> d;
optional<pair<F, C>> step(int S, int T) {
                                                             4.10 Min Cost Circulation [ea0477]
  queue<int> q;
                                                            template <typename F, typename C>
  for (q.push(S), d[S] = 0, up[S] = INF_F;
                                                            struct MinCostCirculation {
    not q.empty(); q.pop()) {
                                                             struct ep { int to; F flow; C cost; };
   int u = q.front(); inq[u] = false;
                                                             int n; vector<int> vis; int visc;
  if (up[u] == 0) continue;
for (int i = 0; i < int(g[u].size()); ++i) {</pre>
                                                             vector<int> fa, fae; vector<vector<int>> g;
                                                             vector<ep> e; vector<C> pi;
    auto e = g[u][i]; int v = e.to;
                                                             MinCostCirculation(int n_) : n(n_), vis(n), visc(0), g
    if (e.f <= 0 or d[v] <= d[u] + e.c) continue;</pre>
                                                                 (n), pi(n) {}
    d[v] = d[u] + e.c; f[v] = \{u, i\};
                                                             void add_edge(int u, int v, F fl, C cs) {
    up[v] = min(up[u], e.f);
                                                              g[u].emplace_back((int)e.size());
    if (not inq[v]) q.push(v);
                                                              e.emplace_back(v, fl, cs);
    inq[v] = true;
                                                              g[v].emplace_back((int)e.size());
  }
                                                              e.emplace_back(u, 0, -cs);
  if (d[T] == INF_C) return nullopt;
                                                             C phi(int x) {
  for (int i = T; i != S; i = f[i].first) {
                                                              if (fa[x] == -1) return 0;
   auto &eg = g[f[i].first][f[i].second];
                                                              if (vis[x] == visc) return pi[x];
   eg.f -= up[T]; g[eg.to][eg.r].f += up[T];
                                                              vis[x] = visc;
 }
                                                              return pi[x] = phi(fa[x]) - e[fae[x]].cost;
  return pair{up[T], d[T]};
                                                             int lca(int u, int v) {
public:
                                                              for (; u != -1 || v != -1; swap(u, v)) if (u != -1) {
   if (vis[u] == visc) return u;
MCMF(int n) : g(n), f(n), inq(n), up(n), d(n, INF_C) {}
void add_edge(int s, int t, F c, C w) {
                                                               vis[u] = visc; u = fa[u];
  g[s].emplace_back(t, int(g[t].size()), c, w);
  g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
                                                              return -1;
pair<F, C> solve(int a, int b) {
                                                             void pushflow(int x, C &cost) {
 F c = 0; C w = 0;
                                                              int v = e[x ^ 1].to, u = e[x].to; ++visc;
 while (auto r = step(a, b)) {
                                                              if (int w = lca(u, v); w == −1) {
   c += r->first, w += r->first * r->second;
                                                               while (v != -1)
   ranges::fill(inq, false); ranges::fill(d, INF_C);
                                                                swap(x ^= 1, fae[v]), swap(u, fa[v]), swap(u, v);
                                                              } else {
  return {c, w};
                                                                int z = u, dir = 0; F f = e[x].flow;
}
                                                                vector<int> cyc = {x};
};
                                                                for (int d : {0, 1})
4.9
      Dijkstra Cost Flow [docfd9]
                                                                for (int i = (d ? u : v); i != w; i = fa[i]) {
template <typename F, typename C> class MCMF {
   static constexpr F INF_F = numeric_limits<F>::max();
                                                                  cyc.push_back(fae[i] ^ d);
                                                                  if (chmin(f, e[fae[i] ^ d].flow)) z = i, dir = d;
static constexpr C INF_C = numeric_limits<C>::max();
                                                               for (int i : cyc) {
struct E { int to, r; F f; C c; };
                                                                e[i].flow -= f; e[i ^ 1].flow += f;
 vector<vector<E>> g; vector<pair<int, int>> f;
vector<F> up; vector<C> d, h;
                                                                cost += f * e[i].cost;
optional<pair<F, C>> step(int S, int T) {
                                                               if (dir) x ^= 1, swap(u, v);
 priority_queue<pair<C, int>> q;
                                                               while (u != z)
  q.emplace(d[S] = 0, S), up[S] = INF_F;
 while (not q.empty()) {
                                                                swap(x ^= 1, fae[v]), swap(u, fa[v]), swap(u, v);
   auto [l, u] = q.top(); q.pop();
   if (up[u] == 0 or l != -d[u]) continue;
   for (int i = 0; i < int(g[u].size()); ++i) {</pre>
                                                             void dfs(int u) {
    auto e = g[u][i]; int v = e.to;
                                                              vis[u] = visc;
    auto nd = d[u] + e.c + h[u] - h[v];
                                                              for (int i : g[u])
    if (e.f <= 0 or d[v] <= nd) continue;</pre>
                                                               if (int v = e[i].to; vis[v] != visc and e[i].flow)
    f[v] = {u, i}; up[v] = min(up[u], e.f);
                                                                fa[v] = u, fae[v] = i, dfs(v);
    q.emplace(-(d[v] = nd), v);
                                                             C simplex() {
```

```
fa.assign(g.size(), -1); fae.assign(g.size(), -1);
  C cost = 0; ++visc; dfs(0);
  for (int fail = 0; fail < ssize(e); )</pre>
   for (int i = 0; i < ssize(e); i++)
if (e[i].flow and e[i].cost < phi(e[i ^ 1].to) -</pre>
    phi(e[i].to))
     fail = 0, pushflow(i, cost), ++visc;
    else ++fail;
  return cost;
 }
4.11 General Matching [5f2293]
struct Matching {
 queue<int> q; int ans, n;
 vector<int> fa, s, v, pre, match;
 int Find(int u) {
  return u == fa[u] ? u : fa[u] = Find(fa[u]); }
 int LCA(int x, int y) {
  static int tk = 0; tk++; x = Find(x); y = Find(y);
  for (;; swap(x, y)) if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
  }
 void Blossom(int x, int y, int l) {
for (; Find(x) != l; x = pre[y]) {
   pre[x] = y, y = match[x];
   if (s[y] == 1) q.push(y), s[y] = 0;
   for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
  }
 bool Bfs(auto &&g, int r) {
  iota(all(fa), 0); ranges::fill(s, -1);
  q = queue<int>(); q.push(r); s[r] = 0;
  for (; !q.empty(); q.pop()) {
   for (int x = q.front(); int u : g[x])
    if (s[u] == -1) {
     if (pre[u] = x, s[u] = 1, match[u] == n) {
      for (int a = u, b = x, last;
         b != n; a = last, b = pre[a])
        last = match[b], match[b] = a, match[a] = b;
      return true;
     q.push(match[u]); s[match[u]] = 0;
    } else if (!s[u] && Find(u) != Find(x)) {
     int l = LCA(u, x);
     Blossom(x, u, l); Blossom(u, x, l);
  }
  return false;
 \label{eq:matching} \textit{Matching}(\textbf{auto} \ \&\&g) \ : \ \textit{ans}(0), \ \textit{n}(\textbf{int}(\texttt{g.size}())),
 fa(n+1), s(n+1), v(n+1), pre(n+1, n), match(n+1, n) {
  for (int x = 0; x < n; ++x)
   if (match[x] == n) ans += Bfs(g, x);
 } // match[x] == n means not matched
}; // test @ yosupo judge
4.12 Weighted Matching [900530]
#define pb emplace_back
#define REP(i, l, r) for (int i=(l); i<=(r); ++i)
struct WeightGraph { // 1-based</pre>
 static const int inf = INT_MAX;
 struct edge { int u, v, w; }; int n, nx;
 vector<int> lab; vector<vector<edge>> g;
 vector<int> slack, match, st, pa, S, vis;
 vector<vector<int>> flo, flo_from; queue<int> q;
 WeightGraph(int n_): n(n_), nx(n * 2), lab(nx + 1), g(nx + 1, vector<edge>(nx + 1)), slack(nx + 1),
  flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
  match = st = pa = S = vis = slack;
  REP(u, 1, n) REP(v, 1, n) g[u][v] = \{u, v, 0\};
 int ED(edge e) {
  return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
 void update_slack(int u, int x, int &s) {
  if (!s || ED(g[u][x]) < ED(g[s][x])) s = u; }</pre>
 void set_slack(int x) {
  slack[x] = 0;
  REP(u, 1, n)
   if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
    update_slack(u, x, slack[x]);
```

```
void q_push(int x) {
if (x \le n) q.push(x);
else for (int y : flo[x]) q_push(y);
void set_st(int x, int b) {
st[x] = b;
if (x > n) for (int y : flo[x]) set_st(y, b);
vector<int> split_flo(auto &f, int xr) {
auto it = find(all(f), xr);
if (auto pr = it - f.begin(); pr % 2 == 1)
 reverse(1 + all(f)), it = f.end() - pr;
auto res = vector(f.begin(), it);
return f.erase(f.begin(), it), res;
void set_match(int u, int v) {
match[u] = g[u][v].v;
if (u <= n) return;</pre>
int xr = flo_from[u][g[u][v].u];
auto &f = flo[u], z = split_flo(f, xr);
REP(i, 0, int(z.size())-1) set_match(z[i], z[i ^ 1]);
set_match(xr, v); f.insert(f.end(), all(z));
void augment(int u, int v) {
for (;;) {
 int xnv = st[match[u]]; set_match(u, v);
 if (!xnv) return;
 set_match(v = xnv, u = st[pa[xnv]]);
int lca(int u, int v) {
static int t = 0; ++t;
for (++t; u || v; swap(u, v)) if (u) {
 if (vis[u] == t) return u;
  vis[u] = t; u = st[match[u]];
 if (u) u = st[pa[u]];
return 0;
void add_blossom(int u, int o, int v) {
int b = int(find(n + 1 + all(st), 0) - begin(st));
lab[b] = 0, S[b] = 0; match[b] = match[o];
vector<int> f = {o};
for (int x : {u, v}) {
 for (int y; x != o; x = st[pa[y]])
   f.pb(x), f.pb(y = st[match[x]]), q_push(y);
 reverse(1 + all(f));
flo[b] = f; set_st(b, b);
REP(x, 1, nx) g[b][x].w = g[x][b].w = 0;
REP(x, 1, n) flo_from[b][x] = 0;
for (int xs : flo[b]) {
 REP(x, 1, nx)
   if (g[b][x].w == 0 || ED(g[xs][x]) < ED(g[b][x]))</pre>
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
 REP(x, 1, n)
  if (flo_from[xs][x]) flo_from[b][x] = xs;
set_slack(b);
void expand_blossom(int b) {
for (int x : flo[b]) set_st(x, x);
int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
for (int x : split_flo(flo[b], xr)) {
 if (xs == -1) { xs = x; continue; }
  pa[xs] = g[x][xs].u; S[xs] = 1, S[x] = 0;
 slack[xs] = 0; set_slack(x); q_push(x); xs = -1;
for (int x : flo[b])
 if (x == xr) S[x] = 1, pa[x] = pa[b];
 else S[x] = -1, set_slack(x);
st[b] = 0;
bool on_found_edge(const edge &e) {
if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
 int nu = st[match[v]]; pa[v] = e.u; S[v] = 1;
 slack[v] = slack[nu] = 0; S[nu] = 0; q_push(nu);
} else if (S[v] == 0) {
 if (int o = lca(u, v)) add_blossom(u, o, v);
 else return augment(u, v), augment(v, u), true;
```

• $S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$ • $S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$

```
return false;
                                                                       S_2(i,k) = \frac{i!}{k!} [x^i] (e^x - 1)^k
                                                                      Derivatives/Integrals
 bool matching() {
                                                                      Integration by parts: \int_a^b f(x)g(x)dx = [F(x)g(x)]_a^b - \int_a^b F(x)g'(x)dx
                                                                      \frac{d}{dx}\sin^{-1}x = \frac{1}{\sqrt{1-x^2}} \left| \frac{d}{dx}\cos^{-1}x = -\frac{1}{\sqrt{1-x^2}} \right|
                                                                                                                  \frac{d}{1-\tan^{-1}x} = \frac{1}{1+x^2}
  ranges::fill(S, -1); ranges::fill(slack, 0);
                                                                                                         \frac{1}{\sqrt{1-x^2}} \left| \frac{d}{dx} \right|
  q = queue<int>();
                                                                       \frac{d}{dx}\tan x = 1 + \tan^2 x \left| \int_{-\infty}^{\infty} \tan ax = -\frac{\ln|\cos ax|}{a} \right|
  REP(x, 1, nx) if (st[x] == x \&\& !match[x])
   pa[x] = 0, S[x] = 0, q_push(x);
                                                                            \int_{a}^{x^{2}} \frac{\sqrt{\pi}}{2} \operatorname{erf}(x) \left| \int_{a}^{y} x e^{ax} dx = \frac{e^{ax}}{a^{2}} (ax - 1) \right|
  if (q.empty()) return false;
  for (;;) {
                                                                          \sqrt{a^2+x^2}=rac{1}{2}\left(x\sqrt{a^2+x^2}+a^2 \operatorname{asinh}(x/a)
ight)
   while (q.size()) {
    int u = q.front(); q.pop();
                                                                      Extended Euler
    if (S[st[u]] == 1) continue;
                                                                      a^b \equiv \begin{cases} a^{(b \mod \varphi(m)) + \varphi(m)} & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m
    REP(v, 1, n)
      if (g[u][v].w > 0 && st[u] != st[v]) {
                                                                      Pentagonal Number Theorem
       if (ED(g[u][v]) != 0)
                                                                      \prod_{n=1}^{\infty} (1 - x^n) = \sum_{k=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2} = (\sum p(n)x^n)^{-1}
        update_slack(u, st[v], slack[st[v]]);
                                                                      5.3 Integer Division* [cd017d]
       else if (on_found_edge(g[u][v])) return true;
                                                                      lld fdiv(lld a, lld b)
                                                                      { return a / b - (a % b && (a < 0) ^ (b < 0)); }
   }
                                                                      lld cdiv(lld a, lld b)
   int d = inf;
   REP(b, n + 1, nx) if (st[b] == b \&\& S[b] == 1)
                                                                      { return a / b + (a % b && (a < 0) ^ (b > 0)); }
                                                                      5.4 FloorSum [fb5917]
    d = min(d, lab[b] / 2);
                                                                     // @param n `n < 2<sup>3</sup>2`
// @param m `1 <= m < 2<sup>3</sup>2`
   REP(x, 1, nx)
    if (int s = slack[x]; st[x] == x && s && S[x] <= 0)</pre>
                                                                      // @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
      d = min(d, ED(g[s][x]) / (S[x] + 2));
   REP(u, 1, n)
                                                                      llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
    if (S[st[u]] == 1) lab[u] += d;
                                                                       llu ans = 0:
    else if (S[st[u]] == 0) {
                                                                       while (true) {
      if (lab[u] <= d) return false;</pre>
                                                                        if (a >= m) ans += n*(n-1)/2 * (a/m), a %= m;
                                                                        if (b >= m) ans += n * (b/m), b %= m;
      lab[u] -= d;
                                                                        if (llu y_max = a * n + b; y_max >= m) {
   REP(b, n + 1, nx) if (st[b] == b && S[b] >= 0) lab[b] += d * (2 - 4 * S[b]);
                                                                         n = (llu)(y_max / m), b = (llu)(y_max % m);
                                                                         swap(m, a);
   REP(x, 1, nx)
                                                                        } else break;
    if (int s = slack[x]; st[x] == x &&
                                                                       }
       s \&\& st[s] != x \&\& ED(g[s][x]) == 0)
                                                                       return ans;
      if (on_found_edge(g[s][x])) return true;
                                                                      lld floor_sum(lld n, lld m, lld a, lld b) {
   REP(b, n + 1, nx)
    if (st[b] == b && S[b] == 1 && lab[b] == 0)
                                                                       if (a < 0) {
      expand_blossom(b);
                                                                        llu a2 = (a \% m + m), d = (a2 - a) / m;
  return false;
                                                                        ans -= 1ULL * n * (n - 1) / 2 * d; a = a2;
 pair<lld, int> solve() {
                                                                       if (b < 0) {
  ranges::fill(match, 0);
                                                                        llu b2 = (b \% m + m), d = (b2 - b) / m;
                                                                        ans -= 1ULL * n * d; b = b2;
  REP(u, 0, n) st[u] = u, flo[u].clear();
  int w_max = 0;
  REP(u, 1, n) REP(v, 1, n) {
                                                                       return ans + floor_sum_unsigned(n, m, a, b);
   flo_from[u][v] = (u == v ? u : 0);
                                                                      5.5 ModMin [2c021c]
   w_max = max(w_max, g[u][v].w);
                                                                      // min{k | l <= ((ak) mod m) <= r}
  REP(u, 1, n) lab[u] = w_max;
                                                                      optional<llu> mod_min(u32 a, u32 m, u32 l, u32 r) {
  int n_matches = 0; lld tot_weight = 0;
                                                                       if (a == 0) return l ? nullopt : optional{0};
  while (matching()) ++n_matches;
                                                                       if (auto k = llu(l + a - 1) / a; k * a <= r)
                                                                        return k;
  REP(u, 1, n) if (match[u] && match[u] < u)
   tot_weight += g[u][match[u]].w;
                                                                       auto b = m / a, c = m % a;
  return make_pair(tot_weight, n_matches);
                                                                       if (auto y = mod_min(c, a, a - r % a, a - l % a))
                                                                        return (l + *y * c + a - 1) / a + *y * b;
                                                                       return nullopt;
 void set_edge(int u, int v, int w) {
  g[u][v].w = g[v][u].w = w; }
                                                                      5.6 Floor Monoid Product [416e89]
      Math
                                                                      /* template <typename T>
     Common Bounds
                                                                      T brute(llu a, llu b, llu c, llu n, T U, T R) {
 n |234567892050100 n |1001e31e61e91e121e15
for (llu i = 1, l = 0; i <= n; i++, res = res * R)
                                                                        for (llu \ r = (a*i+b)/c; \ l < r; ++l) \ res = res * U;
                                      10 11 12 13 14
\binom{2n}{n} 2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8
                                                                       return res;
 n 2345678
                             10
                                  11 12 13
B<sub>n</sub> 2 5 15 52 203 877 4140 21147 115975 7e5 4e6 3e7
                                                                      template <typename T>
                                                                      T euclid(llu a, llu b, llu c, llu n, T U, T R) {
5.2 Equations
                                                                       if (!n) return T{};
Stirling Number of the First Kind
                                                                       if (b >= c)
S_1(n,k) counts the number of permutations of n elements with k disjoint
                                                                        return mpow(U, b / c) * euclid(a, b % c, c, n, U, R);
cycles
                                                                       if (a >= c)
• S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)
                                                                        return euclid(a % c, b, c, n, U, mpow(U, a / c) \star R);
• S_1(n,i) = [x^i] \left(\prod_{i=0}^{n-1} (x+i)\right), use D&Q and taylor shift.
                                                                       llu m = (u128(a) * n + b) / c;
• S_1(i,k) = \frac{i!}{k!} \left[ x^i \right] \left( \sum_{j \ge 1} \frac{x^j}{j} \right)^k
                                                                       if (!m) return mpow(R, n);
return mpow(R, (c - b - 1) / a) * U
Stirling Number of the Second Kind
                                                                        * euclid(c, (c - b - 1) % a, a, m - 1, R, U)
S_2(n,k) counts the number of ways to partition a set of n elements into k
                                                                        * mpow(R, n - (u128(c) * m - b - 1) / a);
nonempty sets.
```

// time complexity is O(log max(a, b, c))

```
// UUUU R UUUUU R ... UUU R 共 N 個 R, 最後一個必是 R
                                                              vector<P> v(sz);
                                                              for (size_t i = 0; i < a.size(); ++i) v[i].RE(a[i]);</pre>
  一直到第 k 個 R 前總共有 (ak+b)/c 個 U
5.7 ax+by=gcd [d0cbdd]
                                                              for (size_t i = 0; i < b.size(); ++i) v[i].IM(b[i]);</pre>
                                                              fft(v.data(), sz, /*inv=*/false);
auto rev = v; reverse(1 + all(rev));
// ax+ny = 1, ax+ny == ax == 1 \ (mod \ n)
void exgcd(lld x, lld y, lld &g, lld &a, lld &b) {
                                                              for (int i = 0; i < sz; ++i) {</pre>
if (y == 0) g = x, a = 1, b = 0;
                                                              P A = (v[i] + conj(rev[i])) / P(2, 0);
else exgcd(y, x \% y, g, b, a), b -= (x / y) * a;
                                                              P B = (v[i] - conj(rev[i])) / P(0, 2);
                                                               v[i] = A * B;
5.8 Chinese Remainder [d69e74]
// please ensure r_i\in[0,m_i)
                                                              VL c(sz); fft(v.data(), sz, /*inv=*/true);
bool crt(lld &m1, lld &r1, lld m2, lld r2) {
                                                              for (int i = 0; i < sz; ++i) c[i] = roundl(RE(v[i]));</pre>
  if (m2 > m1) swap(m1, m2), swap(r1, r2);
                                                              return c;
 lld g, a, b; exgcd(m1, m2, g, a, b);
  if ((r2 - r1) % g != 0) return false;
                                                             VI convolution_mod(const VI &a, const VI &b) {
 m2 /= g; lld D = (r2 - r1) / g % <math>m2 * a % m2;
                                                              if (a.empty() || b.empty()) return {};
 r1 += (D < 0 ? D + m2 : D) * m1; m1 *= m2;
                                                              const int sz = bit_ceil(a.size() + b.size() - 1);
 assert (r1 >= 0 \&\& r1 < m1);
                                                              vector<P> fa(sz), fb(sz);
  return true;
                                                              for (size_t i = 0; i < a.size(); ++i)</pre>
                                                               fa[i] = P(a[i] & ((1 << 15) - 1), a[i] >> 15);
5.9 DiscreteLog [86e463]
                                                              for (size_t i = 0; i < b.size(); ++i)</pre>
template<typename Int>
                                                               fb[i] = P(b[i] & ((1 << 15) - 1), b[i] >> 15);
Int BSGS(Int x, Int y, Int M) {
                                                              fft(fa.data(), sz); fft(fb.data(), sz);
// x^? \setminus equiv y \pmod{M}
                                                              auto rfa = fa; reverse(1 + all(rfa));
Int t = 1, c = 0, g = 1;
                                                              for (int i = 0; i < sz; ++i) fa[i] *= fb[i];</pre>
for (Int M_{-} = M; M_{-} > 0; M_{-} >>= 1) g = g * x % M;
                                                              for (int i = 0; i < sz; ++i) fb[i] *= conj(rfa[i]);</pre>
 for (g = gcd(g, M); t % g != 0; ++c) {
                                                              fft(fa.data(), sz, true); fft(fb.data(), sz, true);
 if (t == y) return c;
                                                              vector<int> res(sz);
 t = t * x % M;
                                                              for (int i = 0; i < sz; ++i) {</pre>
                                                               lld A = (lld)roundl(RE((fa[i] + fb[i]) / P(2, 0)));
if (y % g != 0) return -1;
                                                               lld C = (lld) roundl(IM((fa[i] - fb[i]) / P(0, 2)));
t /= g, y /= g, M /= g;
                                                               lld B = (lld)roundl(IM(fa[i])); B %= p; C %= p;
Int h = 0, gs = 1;

for (; h * h < M; ++h) gs = gs * x % M;
                                                               res[i] = (A + (B << 15) + (C << 30)) % p;
unordered_map<Int, Int> bs;
                                                              return res;
for (Int s = 0; s < h; bs[y] = ++s) y = y * x % M;
                                                            } // test @ yosupo judge with long double
for (Int s = 0; s < M; s += h) {</pre>
 t = t * gs % M;
                                                             5.13
                                                                   CRT for arbitrary mod [e4dde7]
                                                             const int mod = 1000000007;
 if (bs.count(t)) return c + s + h - bs[t];
                                                             const int M1 = 985661441; // G = 3 for M1, M2, M3
                                                             const int M2 = 998244353;
return -1;
                                                             const int M3 = 1004535809;
5.10 Quadratic Residue [f0baec]
                                                             int superBigCRT(lld A, lld B, lld C) {
                                                             static_assert (M1 < M2 && M2 < M3);</pre>
int get_root(int n, int P) { // ensure 0 <= n < p</pre>
if (P == 2 or n == 0) return n;
                                                              constexpr lld r12 = modpow(M1, M2-2, M2);
                                                              constexpr lld r13 = modpow(M1, M3-2, M3);
auto check = [&](lld x) {
                                                              constexpr lld r23 = modpow(M2, M3-2, M3);
  return modpow(int(x), (P - 1) / 2, P); };
if (check(n) != 1) return -1;
                                                              constexpr lld M1M2 = 1LL * M1 * M2 % mod;
                                                              B = (B - A + M2) * r12 % M2;
mt19937 \text{ rnd}(7122); lld z = 1, w;
                                                              C = (C - A + M3) * r13 % M3;
while (check(w = (z * z - n + P) \% P) != P - 1)
                                                              C = (C - B + M3) * r23 % M3;
 z = rnd() \% P;
                                                              return (A + B * M1 + C * M1M2) % mod;
const auto M = [P, w](auto &u, auto &v) {
 auto [a, b] = u; auto [c, d] = v;
 return make_pair((a * c + b * d % P * w) % P,
                                                             5.14 NTT / FFT [2ac7d2]
    (a * d + b * c) % P);
                                                             template <int mod, int G, int maxn> struct NTT {
};
                                                              static_assert(maxn == (maxn & -maxn));
pair<lld, lld> r(1, 0), e(z, 1);
                                                              int roots[maxn];
for (int q = (P + 1) / 2; q; q >>= 1, e = M(e, e))
                                                              NTT () {
 if (q & 1) r = M(r, e);
                                                               int r = modpow(G, (mod - 1) / maxn);
 return int(r.first); // sqrt(n) mod P where P is prime
                                                               for (int i = maxn >> 1; i; i >>= 1) {
                                                                roots[i] = 1;
5.11 FWT [88a937]
                                                                for (int j = 1; j < i; j++)</pre>
/* or convolution:
                                                                roots[i + j] = mul(roots[i + j - 1], r);
* x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                                r = mul(r, r);
* and convolution:
                                                                // for (int j = 0; j < i; j++) // FFT (tested)
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
                                                                // roots[i+j] = polar<llf>(1, PI * j / i);
void fwt(int x[], int N, bool inv = false) {
for (int d = 1; d < N; d <<= 1)</pre>
  for (int s = 0; s < N; s += d * 2)
                                                              // n must be 2^k, and 0 \le F[i] \le mod
  for (int i = s; i < s + d; i++) {
  int j = i + d, ta = x[i], tb = x[j];</pre>
                                                              void operator()(int F[], int n, bool inv = false) {
                                                              for (int i = 0, j = 0; i < n; i++) {
  if (i < j) swap(F[i], F[j]);</pre>
    x[i] = add(ta, tb); x[j] = sub(ta, tb);
                                                                for (int k = n>>1; (j^=k) < k; k>>=1);
if (!inv) return;
 const int invn = modinv(N);
                                                               for (int s = 1; s < n; s *= 2)
 for (int i = 0; i < N; i++) x[i] = mul(x[i], invn);</pre>
                                                                for (int i = 0; i < n; i += s * 2)
                                                                 for (int j = 0; j < s; j++) {
5.12 Packed FFT [Oa6af5]
                                                                  int a = F[i+j], b = mul(F[i+j+s], roots[s+j]);
VL convolution(const VI &a, const VI &b) {
                                                                  F[i+j] = add(a, b); F[i+j+s] = sub(a, b);
if (a.empty() || b.empty()) return {};
const int sz = bit_ceil(a.size() + b.size() - 1);
                                                               if (!inv) return;
// Should be able to handle N <= 10^5, C <= 10^4
                                                               const int invn = modinv(n);
```

```
for (int i = 0; i < n; i++) F[i] = mul(F[i], invn);</pre>
  reverse(F + 1, F + n);
}
      Formal Power Series [c6b99a]
#define fi(l, r) for (size_t i = (l); i < (r); i++)
using S = vector<int>;
auto Mul(auto a, auto b, size_t sz) {
a.resize(sz), b.resize(sz);
ntt(a.data(), sz); ntt(b.data(), sz);
fi(0, sz) a[i] = mul(a[i], b[i]);
return ntt(a.data(), sz, true), a;
S Newton(const S &v, int init, auto &&iter) {
S Q = { init };
 for (int sz = 2; Q.size() < v.size(); sz *= 2) {</pre>
 S A{begin(v), begin(v) + min(sz, int(v.size()))};
A.resize(sz * 2), Q.resize(sz * 2);
                                                              S Q = Mul(Inv(X), Y);
 iter(Q, A, sz * 2); Q.resize(sz);
return Q.resize(v.size()), Q;
S Inv(const S &v) { // v[0] != 0
return Newton(v, modinv(v[0]),
  [](S &X, S &A, int sz) {
                                                              return {Q, Y};
  ntt(X.data(), sz), ntt(A.data(), sz);
   for (int i = 0; i < sz; i++)</pre>
   X[i] = mul(X[i], sub(2, mul(X[i], A[i])));
   ntt(X.data(), sz, true); });
S Dx(S A) {
fi(1, A.size()) A[i - 1] = mul(i, A[i]);
return A.empty() ? A : (A.pop_back(), A);
S Sx(S A) {
A.insert(A.begin(), 0);
fi(1, A.size()) A[i] = mul(modinv(int(i)), A[i]);
return A;
S Ln(const S &A) { // coef[0] == 1; res[0] == 0
auto B = Sx(Mul(Dx(A), Inv(A), bit_ceil(A.size()*2)));
                                                              } // Bostan-Mori
return B.resize(A.size()), B;
S Exp(const S &v) { // coef[0] == 0; res[0] == 1
                                                             ans[0] = tmp[0] = 1;
return Newton(v, 1,
  [](S &X, S &A, int sz) {
   auto Y = X; Y.resize(sz / 2); Y = Ln(Y);
   fi(0, Y.size()) Y[i] = sub(A[i], Y[i]);
   Y[0] = add(Y[0], 1); X = Mul(X, Y, sz); \});
S Pow(S a, lld M) { // period mod*(mod-1)
assert(!a.empty() && a[0] != 0);
const auto imul = [&a](int s) {
  for (int &x: a) x = mul(x, s); }; int c = a[0];
imul(modinv(c)); a = Ln(a); imul(int(M % mod));
a = Exp(a); imul(modpow(c, int(M % (mod - 1))));
                                                              if (n <= 1) return 0;
return a; // mod x^N where N=a.size()
S Sqrt(const S &v) { // need: QuadraticResidue
assert(!v.empty() && v[0] != 0);
const int r = get_root(v[0]); assert(r != -1);
return Newton(v, r,
 [](S &X, S &A, int sz) {
   auto Y = X; Y.resize(sz / 2);
   auto B = Mul(A, Inv(Y), sz);
   for (int i = 0, inv2 = mod / 2 + 1; i < sz; i++)</pre>
                                                               skip[p] = 1;
   X[i] = mul(inv2, add(X[i], B[i])); });
                                                               int ns = 0;
S Mul(auto &&a, auto &&b) {
const auto n = a.size() + b.size() - 1;
auto R = Mul(a, b, bit_ceil(n));
return R.resize(n), R;
S MulT(S a, S b, size_t k) {
                                                               }
assert(b.size()); reverse(all(b)); auto R = Mul(a, b);
                                                               z.resize(ns);
R = vector(R.begin() + b.size() - 1, R.end());
return R.resize(k), R;
S Eval(const S &f, const S &x) {
if (f.empty()) return vector(x.size(), 0);
 const int n = int(max(x.size(), f.size()));
auto q = vector(n * 2, S(2, 1)); S ans(n);
```

```
fi(0, x.size()) q[i + n][1] = sub(0, x[i]);
 for (int i = n - 1; i > 0; i--)
  q[i] = Mul(q[i << 1], q[i << 1 | 1]);
 q[1] = MulT(f, Inv(q[1]), n);
 for (int i = 1; i < n; i++) {</pre>
  auto L = q[i << 1], R = q[i << 1 | 1];</pre>
  q[i << 1 | 0] = MulT(q[i], R, L.size());</pre>
  q[i << 1 | 1] = MulT(q[i], L, R.size());
 for (int i = 0; i < n; i++) ans[i] = q[i + n][0];</pre>
 return ans.resize(x.size()), ans;
pair<S, S> DivMod(const S &A, const S &B) {
 assert(!B.empty() && B.back() != 0);
 if (A.size() < B.size()) return {{}}, A};</pre>
 const auto sz = A.size() - B.size() + 1;
S X = B; reverse(all(X)); X.resize(sz);
 S Y = A; reverse(all(Y)); Y.resize(sz);
 Q.resize(sz); reverse(all(Q)); X = Mul(Q, B); Y = A;
 fi(0, Y.size()) Y[i] = sub(Y[i], X[i]);
 while (Y.size() && Y.back() == 0) Y.pop_back();
 while (Q.size() && Q.back() == 0) Q.pop_back();
} // empty means zero polynomial
int LinearRecursionKth(S a, S c, int64_t k) {
 const auto d = a.size(); assert(c.size() == d + 1);
 const auto sz = bit_ceil(2 * d + 1), o = sz / 2;
 S q = c; for (int &x: q) x = sub(0, x); q[0]=1;
 S p = Mul(a, q); p.resize(sz); q.resize(sz);
 for (int r; r = (k & 1), k; k >>= 1) {
  fill(d + all(p), 0); fill(d + 1 + all(q), 0);
  ntt(p.data(), sz); ntt(q.data(), sz);
  for (size_t i = 0; i < sz; i++)</pre>
   p[i] = mul(p[i], q[(i + o) & (sz - 1)]);
  for (size_t i = 0, j = 0; j < sz; i++, j++)</pre>
   q[i] = q[j] = mul(q[i], q[j]);
  ntt(p.data(), sz, true); ntt(q.data(), sz, true);
for (size_t i = 0; i < d; i++) p[i] = p[i << 1 | r];
for (size_t i = 0; i <= d; i++) q[i] = q[i << 1];</pre>
 return mul(p[0], modinv(q[0]));
} // a_n = \sum_{c,j} a_{(n-j)}, c_0 is not used 5.16 Partition Number [9bb845]
for (int i = 1; i * i <= n; i++) {
 for (int rep = 0; rep < 2; rep++)
for (int j = i; j <= n - i * i; j++)</pre>
   modadd(tmp[j], tmp[j-i]);
 for (int j = i * i; j <= n; j++)</pre>
  modadd(ans[j], tmp[j - i * i]);
        Pi Count [715863]
struct S { int rough; lld large; int id; };
lld PrimeCount(lld n) { // n \sim 10^{13} \Rightarrow < 1s
 const int v = static_cast<int>(sqrtl(n)); int pc = 0;
 vector<int> smalls(v + 1), skip(v + 1); vector<S> z;
 for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
 for (int i : views::iota(0, (v + 1) / 2))
z.emplace_back(2*i+1, (n / (2*i+1) + 1) / 2, i);
 for (int p = 3; p <= v; ++p)
  if (smalls[p] > smalls[p - 1]) {
  const int q = p * p; ++pc;
  if (1LL * q * q > n) break;
  for (int i = q; i <= v; i += 2 * p) skip[i] = 1;
  for (auto e : z) if (!skip[e.rough]) {
   lld d = 1LL * e.rough * p;
   e.large += pc - (d <= v ? z[smalls[d] - pc].large :
  smalls[n / d]);</pre>
   e.id = ns; z[ns++] = e;
  for (int j = v / p; j >= p; --j) {
  int c = smalls[j] - pc, e = min(j * p + p, v + 1);
  for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
 lld ans = z[0].large; z.erase(z.begin());
 for (auto &[rough, large, k] : z) {
```

```
const lld m = n / rough; --k;
                                                                do {
  ans -= large - (pc + \bar{k});
                                                                 p = mmul(msub(z = f(f(z)), y = f(y), n), p, n);
                                                                 if (++i &= 63) if (i == (i & -i)) t = gcd(p, n);
  for (auto [p, \_, l] : z)
   if (l >= k || p * p > m) break;
                                                                } while (t == 1);
   else ans += smalls[m / p] - (pc + l);
                                                                return t == n ? pollard_rho(n) : t;
                                                               \} // test @ yosupo judge, ~270ms for Q=100
                                                                // if use montgomery, ~70ms for Q=100
 return ans;
} // test @ yosupo library checker w/ n=1e11, 68ms
                                                               5.21 Montgomery [648fb3]
5.18 Min 25 Sieve* [3695ef]
                                                               struct Mont { // Montgomery multiplication
                                                                constexpr static int W = 64, L = 6;
template <typename U, typename V> struct min25 {
 lld n; int sq;
                                                                llu mod, R1, R2, xinv;
 vector<U> Ss, Sl, Spre; vector<V> Rs, Rl;
                                                                void set_mod(llu _mod) {
                                                                 mod = _mod; assert(mod & 1); xinv = 1;
 Sieve sv; vector<lld> quo;
 U &S(lld d) { return d < sq ? Ss[d] : Sl[n / d]; }
                                                                 for (int j = 0; j < L; j++) xinv *= 2 - xinv * mod;</pre>
                                                                 assert(xinv * mod == 1);
 V &R(lld d) { return d < sq ? Rs[d] : Rl[n / d]; }</pre>
                                                                 const u128 R = (u128(1) << W) % mod;</pre>
 min25(lld n_{-}) : n(n_{-}), sq((int)sqrt(n) + 1),
                                                                 R1 = llu(R); R2 = llu(R*R \% mod);
 Ss(sq), Sl(sq), Spre(sq), Rs(sq), Rl(sq), sv(sq) {
  for (lld i = 1, Q; i <= n; i = n / Q + 1)</pre>
                                                                llu redc(llu a, llu b) const {
   quo.push_back(Q = n / i);
                                                                 u128 T = u128(a) * b, m = -llu(T) * xinv;
                                                                 T += m * mod; T >>= W;
 U F_prime(auto &&f, auto &&F) {
  for (lld p : sv.primes) Spre[p] = f(p);
                                                                 return llu(T >= mod ? T - mod : T);
  for (int i = 1; i < sq; i++) Spre[i] += Spre[i - 1];</pre>
  for (lld i : quo) S(i) = F(i) - F(1);
                                                                llu from(llu x) const {
  for (lld p : sv.primes)
                                                                 assert(x < mod); return redc(x, R2); }</pre>
   for (lld i : quo) {
                                                                llu get(llu a) const { return redc(a, 1); }
    if (p * p > i) break;
                                                                llu one() const { return R1; }
    S(i) = f(p) * (S(i / p) - Spre[p - 1]);
                                                               } mont;
                                                               // a * b % mod == get(redc(from(a), from(b)))
  return S(n);
                                                               5.22 Berlekamp Massey [a94d00]
 } // F_prime: \sum _ {p is prime, p <= n} f(p)
V F_comp(auto &&g) {</pre>
                                                               template <typename T>
                                                               vector<T> BerlekampMassey(const vector<T> &output) {
  for (lld i : quo) R(i) = V(S(i));
                                                                vector<T> d(output.size() + 1), me, he;
  for (lld p : sv.primes | views::reverse)
                                                                for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
   for (lld i : quo) {
                                                                 for (size_t j = 0; j < me.size(); ++j)
d[i] += output[i - j - 2] * me[j];
if ((d[i] -= output[i - 1]) == 0) continue;</pre>
    if (p * p > i) break;
    lld prod = p;
    for (int c = 1; prod * p <= i; ++c, prod *= p) {</pre>
                                                                 if (me.empty()) {
     R(i) += g(p, c) * (R(i / prod) - V(Spre[p]));
                                                                  me.resize(f = i):
     R(i) += g(p, c + 1);
                                                                  continue;
                                                                 vector<T> o(i - f - 1);
  return R(n);
                                                                 T k = -d[i] / d[f]; o.push_back(-k);
 } // F_comp: \sum _ {2 <= i <= n} g(i)
                                                                 for (T x : he) o.push_back(x * k);
}; // O(n^{3/4} / log n)
                                                                 if (o.size() < me.size()) o.resize(me.size());</pre>
/* U, V 都是環,記 h: U -> V 代表 U 轉型成 V 的函數。
                                                                 for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
要求 h(x + y) = h(x) + h(y); f: lld \rightarrow U 是完全積性;
                                                                 if (i-f+he.size() >= me.size()) he = me, f = i;
g 是積性函數且 h(f(p)) = g(p) 對於質數 p。
                                                                 me = o;
呼叫 F_comp 前需要先呼叫 F_prime 得到 S(i)。
                                                                }
S(i), R(i) 是 F_prime 和 F_comp 在 n/k 點的値。 F(i) = \sum_{i=1}^{n} f(i) 和 f(i) 需要快速求値。
                                                                return me;
g(p, c) := g(pow(p, c)) 需要快速求值。
                                                               5.23 Gauss Elimination [fa0977]
例如若 g(p) 是度數 d 的多項式則可以構造 f(p) 是維護
                                                               using VI = vector<int>; // be careful if A.empty()
using VVI = vector<VI>; // ensure that 0 <= x < mod
pair<VI, VVI> gauss(VVI A, VI b) { // solve Ax=b
pow(p, c) 的 (d+1)-tuple */
5.19 Miller Rabin [fbd812]
bool isprime(llu x) {
                                                                const int N = (int)A.size(), M = (int)A[0].size();
 auto witn = [&](llu a, int t) {
                                                                vector<int> depv, free(M, true); int rk = 0;
  for (llu a2; t--; a = a2) {
                                                                for (int i = 0; i < M; i++) {</pre>
   a2 = mmul(a, a, x);
                                                                 int p = -1;
   if (a2 == 1 && a != 1 && a != x - 1) return true;
                                                                 for (int j = rk; j < N; j++)
if (p == -1 || abs(A[j][i]) > abs(A[p][i]))
  }
  return a != 1;
                                                                 p = j;
if (p == -1 || A[p][i] == 0) continue;
 if (x <= 2 || ~x & 1) return x == 2;
                                                                 swap(A[p], A[rk]); swap(b[p], b[rk]);
 int t = countr_zero(x-1); llu odd = (x-1) >> t;
                                                                 const int inv = modinv(A[rk][i]);
 for (llu m:
                                                                 for (int &x : A[rk]) x = mul(x, inv);
  {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
                                                                 b[rk] = mul(b[rk], inv);
for (int j = 0; j < N; j++) if (j != rk) {</pre>
  if (m % x != 0 && witn(mpow(m % x, odd, x), t))
   return false;
                                                                  int z = A[j][i];
 return true;
                                                                   for (int k = 0; k < M; k++)
} // test @ luogu 143 & yosupo judge, ~1700ms for Q=1e5 // if use montgomery, ~250ms for Q=1e5
                                                                   A[j][k] = sub(A[j][k], mul(z, A[rk][k]));
                                                                  b[j] = sub(b[j], mul(z, b[rk]));
5.20 Pollard Rho [57ad88]
// does not work when n is prime or n == 1
                                                                 depv.push_back(i); free[i] = false; ++rk;
// return any non-trivial factor
                                                                for (int i = rk; i < N; i++)
llu pollard_rho(llu n) {
                                                                 if (b[i] != 0) return {{}}, {{}}}; // not consistent
 static mt19937_64 rnd(120821011);
 if (!(n & 1)) return 2;
                                                                VI x(M); VVI h;
                                                                for (int i = 0; i < rk; i++) x[depv[i]] = b[i];
for (int i = 0; i < M; i++) if (free[i]) {</pre>
 llu y = 2, z = y, c = rnd() % n, p = 1, i = 0, t;
 auto f = [&](llu x) {
 return madd(mmul(x, x, n), c, n); };
                                                                 h.emplace_back(M); h.back()[i] = 1;
```

```
for (int j = 0; j < rk; j++)</pre>
                                                                     p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
                                                                    for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];
   h.back()[depv[j]] = sub(0, A[j][i]);
                                                                    q[n] = -1, d[m + 1][n] = 1;
 return {x, h}; // solution = x + span(h[i])
                                                                    int r = 0;
                                                                    for (int i = 1; i < m; ++i)</pre>
5.24
       CharPoly [cd559d]
                                                                     if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
                                                                    if (d[r][n + 1] < -eps) {</pre>
#define rep(x, y, \bar{z}) for (int x=y; x<z; x++)
using VI = vector<int>; using VVI = vector<VI>;
                                                                     pivot(r, n);
void Hessenberg(VVI &H, int N) {
                                                                     if (!phase(1) || d[m + 1][n + 1] < -eps)
 for (int i = 0; i < N - 2; ++i) {
  for (int j = i + 1; j < N; ++j) if (H[j][i]) {</pre>
                                                                      return VD(n, -inf);
                                                                     for (int i = 0; i < m; ++i) if (p[i] == -1) {
   rep(k, i, N) swap(H[i+1][k], H[j][k]);
                                                                      int s = min_element(d[i].begin(), d[i].end() - 1)
                                                                           - d[i].begin();
   rep(k, 0, N) swap(H[k][i+1], H[k][j]);
   break;
                                                                      pivot(i, s);
  if (!H[i + 1][i]) continue;
  for (int j = i + 2; j < N; ++j) {
                                                                    if (!phase(0)) return VD(n, inf);
   int co = mul(modinv(H[i + 1][i]), H[j][i]);
                                                                    VD x(n);
                                                                    for (int i = 0; i < m; ++i)</pre>
   rep(k, i, N) subeq(H[j][k], mul(H[i+1][k], co));
   rep(k, 0, N) addeq(H[k][i+1], mul(H[k][j], co));
                                                                     if (p[i] < n) x[p[i]] = d[i][n + 1];</pre>
                                                                    return x;
                                                                      // use double instead of long double if possible
}
                                                                   5.26 Simplex Construction
VI CharacteristicPoly(VVI A) {
                                                                   Standard form: maximize \sum_{1 \le i \le n} c_i x_i such that \sum_{1 \le i \le n} A_{ji} x_i \le b_j for all
 int N = (int)A.size(); Hessenberg(A, N);
                                                                   1 \le j \le m and x_i \ge 0 for all 1 \le i \le n.
VVI P(N + 1, VI(N + 1)); P[0][0] = 1;
for (int i = 1; i <= N; ++i) {</pre>
                                                                   1. In case of minimization, let c_i^\prime = -c_i
                                                                   2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} A_{ji} x_i \le -b_j
  rep(j, 0, i+1) P[i][j] = j ? P[i-1][j-1] : 0;
                                                                   3. \sum_{1 \le i \le n}^{-} A_{ji} x_i = b_j \rightarrow \mathsf{add} \le \mathsf{and} \ge 0
  for (int j = i - 1, val = 1; j >= 0; --j) {
  int co = mul(val, A[j][i - 1]);
                                                                   4. If x_i has no lower bound, replace x_i with x_i - x_i'
                                                                   5.27 Adaptive Simpson [b8cef9]
   rep(k, 0, j+1) subeq(P[i][k], mul(P[j][k], co));
                                                                   llf integrate(auto &&f, llf L, llf R) {
   if (j) val = mul(val, A[j][j - 1]);
                                                                    auto simp = [&](llf l, llf r) {
                                                                     llf m = (l + r) / 2;
                                                                     return (f(l) + f(r) + 4.0 * f(m)) * (r - l) / 6.0;
 if (N \& 1) for (int &x: P[N]) x = sub(0, x);
 return P[N]; // test: 2021 PTZ Korea K
                                                                    auto F = [&](auto Y, llf l, llf r, llf v, llf eps) {
                                                                     llf m = (l+r)/2, vl = simp(l, m), vr = simp(m, r);
5.25 Simplex [c9c93b]
                                                                     if (abs(vl + vr - v) <= 15 * eps)
namespace simplex {
                                                                      return vl + vr + (vl + vr - v) / 15.0;
// maximize c^Tx under Ax \le B and x \ge 0
                                                                     return Y(Y, l, m, vl, eps / 2.0) +
/// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
                                                                             Y(Y, m, r, vr, eps / 2.0);
using VD = vector<llf>;
                                                                    return F(F, L, R, simp(L, R), 1e-6);
using VVD = vector<vector<llf>>;
const llf eps = 1e-9, inf = 1e+9;
                                                                   5.28 Poly Roots* [235182]
int n, m; VVD d; vector<int> p, q;
                                                                   VD polyRoots(VD p, llf xmin, llf xmax) {
  if (p.size() == 2) return {-p[0]/p[1]};
void pivot(int r, int s) {
 llf inv = 1.0 / d[r][s];
                                                                    VD d = polyRoots(derivative(p), xmin, xmax), ret;
 for (int i = 0; i < m + 2; ++i)
                                                                    d.pb(xmin-1); d.pb(xmax+1); sort(all(d));
for (size_t i = 0; i + 1 < d.size(); i++) {</pre>
  for (int j = 0; j < n + 2; ++j)</pre>
   if (i != r && j != s)
                                                                     llf l = d[i], h = d[i+1]; bool s = eval(p, l) > 0;
    d[i][j] -= d[r][j] * d[i][s] * inv;
                                                                     if (s ^ (eval(p, h) > 0)) {
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;</pre>
                                                                      for (int _ = 0; _ < 60; _++) {
  llf m = (l + h) / 2, f = eval(p, m);</pre>
 for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
 d[r][s] = inv; swap(p[r], q[s]);
                                                                       ((f \le 0) \land s ? l : h) = m;
bool phase(int z) {
                                                                      ret.push_back((l + h) / 2);
 int x = m + z;
                                                                     }
 while (true) {
  int s = -1;
  for (int i = 0; i <= n; ++i) {
  if (!z && q[i] == -1) continue;</pre>
                                                                    return ret:
                                                                   5.29 Golden Ratio Search* [376bcb]
   if (s == -1 || d[x][i] < d[x][s]) s = i;
                                                                   llf gss(llf a, llf b, auto &&f) {
                                                                    llf r = (sqrt(5)-1)/2, eps = 1e-7;
  if (s == -1 || d[x][s] > -eps) return true;
                                                                    llf x1 = b - r*(b-a), x2 = a + r*(b-a);
  int r = -1;
                                                                    llf f1 = f(x1), f2 = f(x2);
  for (int i = 0; i < m; ++i) {
                                                                    while (b-a > eps)
   if (d[i][s] < eps) continue;</pre>
                                                                     if (f1 < f2) { //change to > to find maximum
   if (r == -1 ||
                                                                      b = x2; x2 = x1; f2 = f1;
    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
                                                                      x1 = b - r*(b-a); f1 = f(x1);
  if (r == -1) return false;
                                                                     } else {
                                                                      a = x1; x1 = x2; f1 = f2;
  pivot(r, s);
                                                                      x2 = a + r*(b-a); f2 = f(x2);
                                                                    return a:
VD solve(const VVD &a, const VD &b, const VD &c) {
m = (int)b.size(), n = (int)c.size();
                                                                        Geometry
 d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)</pre>
                                                                        Basic Geometry [1d2d70]
                                                                   #define IM imag
 for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
 p.resize(m), q.resize(n + 1);
                                                                   #define RE real
 for (int i = 0; i < m; ++i)</pre>
                                                                  using lld = int64_t;
```

```
using llf = long double;
                                                               mx = max(mx, B * sin(deg) * sin(deg));
using PT = complex<lld>;
using PF = complex<llf>;
                                                              return {mn, mx};
using P = PT;
                                                             } // test @ UVA 819
llf abs(P p) { return sqrtl(norm(p)); }
                                                                   Minkowski Sum [602806]
PF toPF(PT p) { return PF{RE(p), IM(p)}; }
                                                             // A, B are strict convex hull rotate to min by (X, Y)
int sgn(lld x) \{ return (x > 0) - (x < 0); \}
                                                             vector<P> Minkowski(vector<P> A, vector<P> B) {
lld dot(P a, P b) { return RE(conj(a) * b); }
                                                              const int N = (int)A.size(), M = (int)B.size();
lld cross(P a, P b) { return IM(conj(a) * b); }
                                                              vector<P> sa(N), sb(M), C(N + M + 1);
for (int i = 0; i < N; i++) sa[i] = A[(i+1)%N]-A[i];</pre>
int ori(P a, P b, P c) {
return sgn(cross(b - a, c - a));
                                                              for (int i = 0; i < M; i++) sb[i] = B[(i+1)%M]-B[i];</pre>
                                                              C[0] = A[0] + B[0];
int quad(P p) {
                                                              for (int i = 0, j = 0; i < N || j < M; ) {
  P e = (j>=M || (i<N && cross(sa[i], sb[j])>=0))
 return (IM(p) == 0) // use sgn for PF
  ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
                                                                ? sa[i++] : sb[j++];
                                                               C[i + j] = e;
int argCmp(P a, P b) {
  // returns 0/+-1, starts from theta = -PI
                                                             partial_sum(all(C), C.begin()); C.pop_back();
return convex_hull(C); // just to remove colinear
} // be careful if min(|A|, |B|)<=2</pre>
 int qa = quad(a), qb = quad(b);
 if (qa != qb) return sgn(qa - qb);
 return sgn(cross(b, a));
                                                             6.6 Segment Intersection [60d016]
                                                             struct Seg { // closed segment
P rot90(P p) { return P{-IM(p), RE(p)}; }
                                                              P st, dir; // represent st + t*dir for 0<=t<=1
template <typename V> llf area(const V & pt) {
                                                              Seg(P s, P e) : st(s), dir(e - s) {}
 lld ret = 0;
                                                              static bool valid(lld p, lld q) {
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
                                                                // is there t s.t. 0 <= t <= 1 && qt == p ?
  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
                                                               if (q < 0) q = -q, p = -p;
 return ret / 2.0;
                                                               return 0 <= p && p <= q;
template <typename V> PF center(const V & pt) {
                                                              vector<P> ends() const { return { st, st + dir }; }
 P ret = 0; lld A = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++) {</pre>
                                                             template <typename T> bool isInter(T A, P p) {
  lld cur = cross(pt[i] - pt[0], pt[i+1] - pt[0]);
                                                              if (A.dir == P(0)) return p == A.st; // BE CAREFUL
  ret += (pt[i] + pt[i + 1] + pt[0]) * cur; A += cur;
                                                              return cross(p - A.st, A.dir) == 0 &&
                                                               T::valid(dot(p - A.st, A.dir), norm(A.dir));
return toPF(ret) / llf(A * 3);
                                                             template <typename U, typename V>
PF project(PF p, PF q) { // p onto q
                                                             bool isInter(U A, V B) {
return dot(p, q) * q / dot(q, q); // dot<llf>
                                                              if (cross(A.dir, B.dir) == 0) { // BE CAREFUL
                                                               bool res = false;
6.2 2D Convex Hull [ecba37]
                                                               for (P p: A.ends()) res |= isInter(B, p);
// from NaCl, counterclockwise, be careful of n<=2
                                                               for (P p: B.ends()) res |= isInter(A, p);
vector<P> convex_hull(vector<P> v) { // n==0 will RE
                                                               return res;
 sort(all(v)); // by X then Y
 if (v[0] == v.back()) return {v[0]};
                                                              P D = B.st - A.st; lld C = cross(A.dir, B.dir);
 int t = 0, s = 1; vector<P> h(v.size() + 1);
                                                              return U::valid(cross(D, B.dir), C) &&
 for (int _ = 2; _--; s = t--, reverse(all(v)))
                                                               V::valid(cross(D, A.dir), C);
  for (P p : v) {
   while (t>s && ori(p, h[t-1], h[t-2]) >= 0) t--;
                                                             6.7 Halfplane Intersection [f2bd8f]
   h[t++] = p;
                                                             struct Line
                                                              P st, ed, dir;
 return h.resize(t), h;
                                                              Line (P s, P e) : st(s), ed(e), dir(e - s) {}
                                                             }; using LN = const Line &;
6.3 2D Farthest Pair [8b5844]
                                                             PF intersect(LN A, LN B) {
// p is CCW convex hull w/o colinear points
                                                              llf t = cross(B.st - A.st, B.dir) /
int n = (int)p.size(), pos = 1; lld ans = 0;
                                                               llf(cross(A.dir, B.dir));
for (int i = 0; i < n; i++) {</pre>
                                                              return toPF(A.st) + toPF(A.dir) * t; // C^3 / C^2
P = p[(i + 1) \% n] - p[i];
 while (cross(e, p[(pos + 1) % n] - p[i]) >
                                                             bool cov(LN l, LN A, LN B) {
     cross(e, p[pos] - p[i]))
                                                              i128 u = cross(B.st-A.st, B.dir);
  pos = (pos + 1) % n;
                                                              i128 v = cross(A.dir, B.dir);
 for (int j: {i, (i + 1) % n})
                                                              // ori(l.st, l.ed, A.st + A.dir*(u/v)) <= 0?
 ans = max(ans, norm(p[pos] - p[j]));
                                                              i128 x = RE(A.dir) * u + RE(A.st - l.st) * v;
} // tested @ AOJ CGL_4_B
                                                              i128 y = IM(A.dir) * u + IM(A.st - l.st) * v;
6.4 MinMax Enclosing Rect [e4470c]
                                                              return sgn(x*IM(l.dir) - y*RE(l.dir)) * sgn(v) >= 0;
// from 8BQube, plz ensure p is strict convex hull
                                                             } // x, y are C^3, also sgn<i128> is needed
const llf INF = 1e18, qi = acos(-1) / 2 * 3;
                                                             bool operator<(LN a, LN b) {</pre>
pair<llf, llf> solve(const vector<P> &p) {
                                                              if (int c = argCmp(a.dir, b.dir)) return c == -1;
 llf mx = 0, mn = INF; int n = (int)p.size();
                                                              return ori(a.st, a.ed, b.st) < 0;</pre>
 for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {
#define Z(v) (p[(v) % n] - p[i])
                                                             // cross(pt-line.st, line.dir)<=0 <-> pt in half plane
  P e = Z(i + 1);
                                                             // the half plane is the LHS when going from st to ed
 while (cross(e, Z(u + 1)) > cross(e, Z(u))) ++u;
                                                             llf HPI(vector<Line> &q) {
  while (dot(e, Z(r + 1)) > dot(e, Z(r))) ++r;
                                                              sort(q.begin(), q.end());
                                                              int n = (int)q.size(), l = 0, r = -1;
  if (!i) l = r + 1;
                                                              for (int i = 0; i < n; i++) {
   if (i && !argCmp(q[i].dir, q[i-1].dir)) continue;</pre>
  while (dot(e, Z(l + 1)) < dot(e, Z(l))) ++l;</pre>
  PD = p[r \% n] - p[l \% n];
  llf H = cross(e, Z(u)) / llf(norm(e));
                                                               while (l < r && cov(q[i], q[r-1], q[r])) --r;</pre>
                                                               while (l < r && cov(q[i], q[l], q[l+1])) ++l;</pre>
  mn = min(mn, dot(e, D) * H);
  llf B = sqrt(norm(D)) * sqrt(norm(Z(u)));
                                                               q[++r] = q[i];
  llf deg = (qi - acos(dot(D, Z(u)) / B)) / 2;
```

bool PIH(const vector<P> &h, P z, bool strict = true) {

```
while (l < r && cov(q[l], q[r-1], q[r])) --r;</pre>
                                                                 int n = (int)h.size(), a = 1, b = n - 1, r = !strict;
                                                                 if (n < 3) return r && isAnti(h[0] - z, h[n-1] - z);</pre>
 while (l < r && cov(q[r], q[l], q[l+1])) ++l;</pre>
 n = r - l + 1; // q[l .. r] are the lines
                                                                 if (ori(h[0],h[a],h[b]) > 0) swap(a, b);
 if (n <= 2 || !argCmp(q[l].dir, q[r].dir)) return 0;</pre>
                                                                 if (ori(h[0],h[a],z) >= r || ori(h[0],h[b],z) <= -r)</pre>
 vector<PF> pt(n);
                                                                 return false;
 for (int i = 0; i < n; i++)</pre>
                                                                 while (abs(a - b) > 1) {
                                                                  int c = (a + b) / 2;
 pt[i] = intersect(q[i+l], q[(i+1)%n+l]);
                                                                  (ori(h[0], h[c], z) > 0 ? b : a) = c;
 return area(pt);
} // test @ 2020 Nordic NCPC : BigBrother
                                                                 return ori(h[a], h[b], z) < r;</pre>
6.8 SegmentDist (Sausage) [9d8603]
// be careful of abs<complex<int>> (replace _abs below)
                                                                      Point In Polygon [037c52]
                                                                6.12
llf PointSegDist(P A, Seg B) {
                                                               bool PIP(const vector<P> &p, P z, bool strict = true) {
 if (B.dir == P(0)) return _abs(A - B.st);
                                                                 int cnt = 0, n = (int)p.size();
 if (sgn(dot(A - B.st, B.dir)) *
                                                                 for (int i = 0; i < n; i++) {</pre>
   sgn(dot(A - B.ed, B.dir)) <= 0)</pre>
                                                                 P A = p[i], B = p[(i + 1) % n];
  return abs(cross(A - B.st, B.dir)) / _abs(B.dir);
                                                                 if (isInter(Seg(A, B), z)) return !strict;
auto zy = IM(z), Ay = IM(A), By = IM(B);
cnt ^= ((zy<Ay) - (zy<By)) * ori(z, A, B) > 0;
 return min(_abs(A - B.st), _abs(A - B.ed));
llf SegSegDist(const Seg &s1, const Seg &s2) {
 if (isInter(s1, s2)) return 0;
                                                                 return cnt;
 return min({
   PointSegDist(s1.st, s2),
                                                                6.13 Point In Polygon (Fast) [2cd3d6]
   PointSegDist(s1.ed, s2),
   PointSegDist(s2.st, s1),
                                                                vector<<mark>int</mark>> PIPfast(vector<P> p, vector<P> q) -
PointSegDist(s2.ed, s1) });
} // test @ Q0J2444 / PTZ19 Summer.D3
                                                                 const int N = int(p.size()), Q = int(q.size());
                                                                 vector<pair<P, int>> evt; vector<Seg> edge;
6.9 Rotating Sweep Line [8aff27]
                                                                 for (int i = 0; i < N; i++) {</pre>
struct Event {
                                                                  int a = i, b = (i + 1) % N;
 P d; int u, v;
                                                                 P A = p[a], B = p[b];
 bool operator<(const Event &b) const {</pre>
                                                                  assert (A < B || B < A); // std::operator<
  return sgn(cross(d, b.d)) > 0; }
                                                                  if (B < A) swap(A, B);
                                                                  evt.emplace_back(A, i); evt.emplace_back(B, ~i);
P makePositive(P z) { return cmpxy(z, 0) ? -z : z; }
                                                                 edge.emplace_back(A, B);
void rotatingSweepLine(const vector<P> &p) {
 const int n = int(p.size());
                                                                 for (int i = 0; i < Q; i++)</pre>
 vector<Event> e; e.reserve(n * (n - 1) / 2);
                                                                 evt.emplace_back(q[i], i + N);
 for (int i = 0; i < n; i++)
for (int j = i + 1; j < n; j++)</pre>
                                                                 sort(all(evt));
                                                                 auto vtx = p; sort(all(vtx));
auto eval = [](const Seg &a, lld x) -> llf {
   e.emplace_back(makePositive(p[i] - p[j]), i, j);
 sort(all(e));
                                                                  if (RE(a.dir) == 0) {
 vector<int> ord(n), pos(n);
                                                                   assert (x == RE(a.st));
 iota(all(ord), 0);
sort(all(ord), [&p](int i, int j) {
                                                                   return IM(a.st) + llf(IM(a.dir)) / 2;
  return cmpxy(p[i], p[j]); });
                                                                  llf t = (x - RE(a.st)) / llf(RE(a.dir));
 for (int i = 0; i < n; i++) pos[ord[i]] = i;
const auto makeReverse = [](auto &v) {</pre>
                                                                  return IM(a.st) + IM(a.dir) * t;
  sort(all(v)); v.erase(unique(all(v)), v.end());
                                                                 lld cur_x = 0;
  vector<pair<int,int>> segs;
for (size_t i = 0, j = 0; i < v.size(); i = j) {</pre>
                                                                 auto cmp = [&](const Seg &a, const Seg &b) -> bool {
                                                                  if (int s = sgn(eval(a, cur_x) - eval(b, cur_x)))
   for (; j < v.size() && v[j] - v[i] <= j - i; j++);</pre>
                                                                   return s == -1; // be careful: sgn<llf>, sgn<lld>
   segs.emplace_back(v[i], v[j - 1] + 1 + 1);
                                                                  int s = sgn(cross(b.dir, a.dir));
                                                                  if (cur_x != RE(a.st) && cur_x != RE(b.st)) s *= -1;
  return segs;
                                                                  return s == -1;
 };
                                                                 };
 for (size_t i = 0, j = 0; i < e.size(); i = j) {</pre>
                                                                 namespace pbds = __gnu_pbds;
 /* do here */
                                                                 pbds::tree<Seg, int, decltype(cmp),</pre>
  vector<size_t> tmp;
                                                                  pbds::rb_tree_tag,
  for (; j < e.size() && !(e[i] < e[j]); j++)</pre>
                                                                  pbds::tree_order_statistics_node_update> st(cmp);
   tmp.push_back(min(pos[e[j].u], pos[e[j].v]));
                                                                 auto answer = [&](P ep) {
  for (auto [l, r] : makeReverse(tmp)) {
                                                                  if (binary_search(all(vtx), ep))
   reverse(ord.begin() + l, ord.begin() + r);
                                                                   return 1; // on vertex
   for (int t = l; t < r; t++) pos[ord[t]] = t;</pre>
                                                                  Seg H(ep, ep); // ??
                                                                  auto it = st.lower_bound(H);
}
                                                                  if (it != st.end() && isInter(it->first, ep))
                                                                   return 1; // on edge
6.10 Hull Cut* [277def]
                                                                  if (it != st.begin() && isInter(prev(it)->first, ep))
vector<P> cut(const vector<P> &p, P s, P e) {
                                                                   return 1; // on edge
 vector<P> res;
                                                                  auto rk = st.order_of_key(H);
 for (size_t i = 0; i < p.size(); i++) {</pre>
                                                                  return rk % 2 == 0 ? 0 : 2; // 0: outside, 2: inside
  P cur = p[i], prv = i ? p[i-1] : p.back();
                                                                 }:
 bool side = ori(s, e, cur) < 0;</pre>
                                                                 vector<int> ans(Q);
  if (side != (ori(s, e, prv) < 0))
                                                                 for (auto [ep, i] : evt) {
   res.push_back(intersect({s, e}, {cur, prv}));
                                                                 cur_x = RE(ep);
 if (side) res.push_back(cur);
                                                                  if (i < 0) { // remove
 } // P is complex<llf>
                                                                   st.erase(edge[~i]);
return res; // hull intersection with halfplane
} // left of the line s -> e
                                                                  } else if (i < N) { // insert</pre>
                                                                   auto [it, succ] = st.insert({edge[i], i});
6.11 Point In Hull [13edeb]
                                                                   assert(succ);
bool isAnti(P a, P b) {
                                                                  } else ans[i - N] = answer(ep);
 return cross(a, b) == 0 && dot(a, b) <= 0; }
```

return ans;

vec = vec * sqrt(r * r - dis * dis) / abs(vec);

```
// test @ AOJ CGL_3_C
                                                              return {ft + vec, ft - vec}; // sqrt_safe?
6.14 Cyclic Ternary Search [162adf]
int cyclic_ternary_search(int N, auto &&lt_) {
                                                             6.19
                                                                    Poly-Circle Intersection [8e5133]
auto lt = [&](int x, int y) {
                                                            // Divides into multiple triangle, and sum up
 return lt_(x % N, y % N); };
                                                             // from 8BQube, test by HDU2892 & AOJ CGL_7_H
int l = 0, r = N; bool up = lt(0, 1);
                                                            llf _area(PF pa, PF pb, llf r) {
                                                              if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
while (r - l > 1) {
 int m = (l + r) / 2;
                                                              if (abs(pb) < eps) return 0;</pre>
 if (lt(m, 0) ? up : !lt(m, m+1)) r = m;
                                                              llf S, h, theta;
 else l = m;
                                                              llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
                                                              llf cB = dot(pb, pb-pa) / a / c, B = acos_safe(cB);
return (lt(l, r) ? r : l) % N;
                                                              llf cC = dot(pa, pb) / a / b, C = acos_safe(cC);
} // find maximum; be careful if N == 0
                                                              if (a > r) {
      Tangent of Points to Hull [8e1343]
                                                               S = (C / 2) * r * r; h = a * b * sin(C) / c;
                                                               if (h < r && B < PI / 2)
pair<int, int> get_tangent(const vector<P> &v, P p) {
                                                                S = (acos\_safe(h/r)*r*r - h*sqrt\_safe(r*r-h*h));
auto gao = [&](int s) {
                                                              } else if (b > r) {
 return cyclic_ternary_search(v.size(),
                                                               theta = PI - B - asin\_safe(sin(B) / r * a);
    [&](int x, int y) {
                                                               S = 0.5 * a*r*sin(theta) + (C-theta)/2 * r * r;
     return ori(p, v[x], v[y]) == s; });
}; // test @ codeforces.com/gym/101201/problem/E
                                                              } else
                                                               S = 0.5 * sin(C) * a * b;
return {gao(1), gao(-1)}; // (a,b):ori(p,v[a],v[b])<0
} // plz ensure that point strictly out of hull
                                                              return S;
    if colinear, returns arbitrary point on line
6.16 Circle Class & Intersection [d5df51]
                                                             llf area_poly_circle(const vector<PF> &v, PF 0, llf r)
llf FMOD(llf x) {
                                                              llf S = 0;
if (x < -PI) x += PI * 2;
                                                              for (size_t i = 0, N = v.size(); i < N; ++i)</pre>
if (x > PI) x -= PI * 2;
                                                               S += _area(v[i] - 0, v[(i + 1) % N] - 0, r) *
return x;
                                                                  ori(0, v[i], v[(i + 1) % N]);
                                                              return abs(S);
struct Cir { PF o; llf r; };
// be carefule when tangent
                                                             6.20 Min Covering Circle [054ee0]
vector<llf> intersectAngle(Cir a, Cir b) {
                                                            Cir getCircum(P a, P b, P c){ // P = complex<llf>
    P z1 = a - b, z2 = a - c; llf D = cross(z1, z2) * 2;
PF dir = b.o - a.o; llf d2 = norm(dir);

if (norm(a.r - b.r) >= d2) { // norm(x) := |x|^2
 if (a.r < b.r) return {-PI, PI}; // a in b</pre>
                                                              auto c1 = dot(a + b, z1), c2 = dot(a + c, z2);
 else return {}; // b in a
                                                              P o = rot90(c2 * z1 - c1 * z2) / D;
} else if (norm(a.r + b.r) <= d2) return {};</pre>
                                                              return { o, abs(o - a) };
llf dis = abs(dir), theta = arg(dir);
llf phi = acos((a.r * a.r + d2 - b.r * b.r) /
                                                             Cir minCircleCover(vector<P> p) { // what if p.empty?
   (2 * a.r * dis)); // is acos_safe needed ?
                                                              Cir c = { 0, 0 }; shuffle(all(p), mt19937(114514));
llf L = FMOD(theta - phi), R = FMOD(theta + phi);
                                                              for (size_t i = 0; i < p.size(); i++) {</pre>
return { L, R };
                                                               if (abs(p[i] - c.o) <= c.r) continue;</pre>
                                                               c = { p[i], 0 };
vector<PF> intersectPoint(Cir a, Cir b) {
                                                               for (size_t j = 0; j < i; j++) {</pre>
llf d = abs(a.o - b.o);
                                                                if (abs(p[j] - c.o) <= c.r) continue;</pre>
if (d > b.r+a.r || d < abs(b.r-a.r)) return {};</pre>
                                                                c.o = (p[i] + p[j]) / llf(2);
llf dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
                                                                c.r = abs(p[i] - c.o);
PF dir = (a.o - b.o) / d;
                                                                for (size_t k = 0; k < j; k++) {</pre>
PF u = dir * d1 + b.o;
                                                                 if (abs(p[k] - c.o) <= c.r) continue;</pre>
PF v = rot90(dir) * sqrt(max(0.0L, b.r*b.r-d1*d1));
                                                                 c = getCircum(p[i], p[j], p[k]);
return \{u + v, u - v\};
} // test @ AOJ CGL probs
                                                               }
6.17
      Circle Common Tangent [d97f1c]
                                                              return c;
// be careful of tangent / exact same circle
                                                             } // test @ TIOJ 1093 & luogu P1742
// sign1 = 1 for outer tang, -1 for inter tang
vector<Line> common_tan(const Cir &a, const Cir &b, int
                                                             6.21 Circle Union [073c1c]
     sign1) {
                                                             #define eb emplace_back
if (norm(a.o - b.o) < eps) return {};</pre>
                                                             struct Teve { // test@SPOJ N=1000, 0.3~0.5s
llf d = abs(a.o - b.o), c = (a.r - sign1 * b.r) / d;
                                                              PF p; llf a; int add; // point, ang, add
PF v = (b.o - a.o) / d;
                                                              Teve(PF x, llf y, int z) : p(x), a(y), add(z) {}
if (c * c > 1) return {};
                                                              bool operator<(Teve &b) const { return a < b.a; }</pre>
if (abs(c * c - 1) < eps) {
 PF p = a.o + c * v * a.r;
                                                             // strict: x = 0, otherwise x = -1
 return {Line(p, p + rot90(b.o - a.o))};
                                                             bool disjunct(Cir &a, Cir &b, int x)
                                                             { return sgn(abs(a.o - b.o) - a.r - b.r) > x; }
vector<Line> ret; llf h = sqrt(max(0.0L, 1-c*c));
for (int sign2 : {1, -1}) {
                                                            bool contain(Cir &a, Cir &b, int x)
                                                             { return sgn(a.r - b.r - abs(a.o - b.o)) > x; }
 PF n = c * v + sign2 * h * rot90(v);
                                                             vector<llf> CircleUnion(vector<Cir>> &c) {
 PF p1 = a.o + n * a.r;
                                                              // area[i] : area covered by at least i circles
 PF p2 = b.o + n * (b.r * sign1);
                                                              int N = (int)c.size(); vector<llf> area(N + 1);
 ret.emplace_back(p1, p2);
                                                              vector<vector<int>> overlap(N, vector<int>(N));
}
                                                              auto g = overlap; // use simple 2darray to speedup
                                                              for (int i = 0; i < N; ++i)</pre>
return ret;
                                                               for (int j = 0; j < N; ++j) {
      Line-Circle Intersection [10786a]
                                                                /* c[j] is non-strictly in c[i]. */
                                                                overlap[i][j] = i != j &&
vector<PF> LineCircleInter(PF p1, PF p2, PF o, llf r) {
                                                                 (sgn(c[i].r - c[j].r) > 0 | |
PF ft = p1 + project(o-p1, p2-p1), vec = p2-p1;
                                                                  (sgn(c[i].r - c[j].r) == 0 \&\& i < j)) \&\&
llf dis = abs(o - ft);
if (abs(dis - r) < eps) return {ft};</pre>
                                                                 contain(c[i], c[j], -1);
if (dis > r) return {};
```

for (int i = 0; i < N; ++i)</pre>

```
for (int j = 0; j < N; ++j)
g[i][j] = i != j && !(overlap[i][j] ||</pre>
                                                                 int a, b, c;
                                                                 Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
      overlap[j][i] \mid\mid disjunct(c[i], c[j], -1));
                                                                };
  for (int i = 0; i < N; ++i) {</pre>
                                                                auto preprocess(const vector<P3> &pt) {
   vector<Teve> eve; int cnt = 1;
                                                                 auto G = pt.begin();
   for (int j = 0; j < N; ++j) cnt += overlap[j][i];</pre>
                                                                 auto a = find_if(all(pt), [&](P3 z) {
   // if (cnt > 1) continue; (if only need area[1])
                                                                  return z != *G; }) - G;
   for (int j = 0; j < N; ++j) if (g[i][j]) {</pre>
                                                                 auto b = find_if(all(pt), [&](P3 z) {
    auto IP = intersectPoint(c[i], c[j]);
                                                                  return ver(*G, pt[a], z) != P3(0, 0, 0); }) - G;
    PF aa = IP[1], bb = IP[0];
                                                                 auto c = find_if(all(pt), [&](P3 z) {
    llf A = arg(aa - c[i].o), B = arg(bb - c[i].o);
                                                                  return volume(*G, pt[a], pt[b], z) != 0; }) - G;
    eve.eb(bb, B, 1); eve.eb(aa, A, -1);
                                                                 vector<size_t> id;
                                                                 for (size_t i = 0; i < pt.size(); i++)</pre>
    if (B > A) ++cnt;
                                                                  if (i != a && i != b && i != c) id.push_back(i);
   if (eve.empty()) area[cnt] += PI*c[i].r*c[i].r;
                                                                 return tuple{a, b, c, id};
   else {
    sort(eve.begin(), eve.end());
                                                                // return the faces with pt indexes
                                                                // all points coplanar case will WA
    eve.eb(eve[0]); eve.back().a += PI * 2;
    for (size_t j = 0; j + 1 < eve.size(); j++) {</pre>
                                                                vector<Face> convex_hull_3D(const vector<P3> &pt) {
     cnt += eve[j].add;
                                                                 const int n = int(pt.size());
                                                                 if (n <= 3) return {}; // be careful about edge case</pre>
     area[cnt] += cross(eve[j].p, eve[j+1].p) \star.5;
     llf t = eve[j + 1].a - eve[j].a;
                                                                 vector<Face> now;
     area[cnt] += (t-sin(t)) * c[i].r * c[i].r *.5;
                                                                 vector<vector<int>> z(n, vector<int>(n));
                                                                 auto [a, b, c, ord] = preprocess(pt);
  }
                                                                 now.emplace_back(a, b, c); now.emplace_back(c, b, a);
                                                                 for (auto i : ord) {
 return area;
                                                                  vector<Face> next;
                                                                  for (const auto &f : now) {
                                                                   lld v = volume(pt[f.a], pt[f.b], pt[f.c], pt[i]);
       Polygon Union [42e75b]
                                                                   if (v <= 0) next.push_back(f);</pre>
 llf polyUnion(const vector<vector<P>> &p) {
                                                                   z[f.a][f.b] = z[f.b][f.c] = z[f.c][f.a] = sgn(v);
 vector<tuple<P, P, int>> seg;
for (int i = 0; i < ssize(p); i++)</pre>
                                                                  const auto F = [\&](int x, int y) \{
  for (int j = 0, m = int(p[i].size()); j < m; j++)</pre>
                                                                   if (z[x][y] > 0 && z[y][x] <= 0)
    seg.emplace_back(p[i][j], p[i][(j + 1) % m], i);
                                                                    next.emplace_back(x, y, i);
 llf ret = 0; // area of p[i] must be non-negative
 for (auto [A, B, i] : seg) {
                                                                  for (const auto &f : now)
  vector<pair<llf, int>> evt{{0, 0}, {1, 0}};
for (auto [C, D, j] : seg) {
                                                                   F(f.a, f.b), F(f.b, f.c), F(f.c, f.a);
                                                                  now = next;
    int sc = ori(A, B, C), sd = ori(A, B, D);
    if (sc != sd && i != j && min(sc, sd) < 0) {</pre>
                                                                 return now;
     llf sa = cross(D-C, A-C), sb = cross(D-C, B-C);
     evt.emplace_back(sa / (sa - sb), sgn(sc - sd));
                                                                // n^2 delaunay: facets with negative z normal of
    } else if (!sc && !sd && j < i
                                                                // convexhull of (x, y, x^2 + y^2), use a pseudo-point
      && sgn(dot(B - A, D - C)) > 0) {
                                                                // (0, 0, inf) to avoid degenerate case
     evt.emplace_back(real((C - A) / (B - A)), 1);
                                                                // test @ SPOJ CH3D
     evt.emplace_back(real((D - A) / (B - A)), -1);
                                                                // llf area = 0, vol = 0; // surface area / volume
   }
                                                                // for (auto [a, b, c]: faces)
                                                                    area += abs(ver(p[a], p[b], p[c]))/2.0,
  for (auto &[q, _] : evt) q = clamp<llf>(q, 0, 1);
                                                                    vol += volume(P3(0, 0, 0), p[a], p[b], p[c])/6.0;
   sort(evt.begin(), evt.end());
                                                                6.25 3D Projection [68f350]
   llf sum = 0, last = 0; int cnt = 0;
                                                                using P3F = valarray<llf>;
   for (auto [q, c] : evt) {
                                                                P3F toP3F(P3 p) { return {p.x, p.y, p.z}; }
   if (!cnt) sum += q - last;
                                                                llf dot(P3F a, P3F b) {
    cnt += c; last = q;
                                                                 return a[0]*b[0]+a[1]*b[1]+a[2]*b[2];
  ret += cross(A, B) * sum;
                                                                P3F housev(P3 A, P3 B, int s) {
                                                                 const llf a = abs(A), b = abs(B);
 return ret / 2;
                                                                 return toP3F(A) / a + s \star toP3F(B) / b;
 6.23 3D Point [46b73b]
                                                                P project(P3 p, P3 q) {
struct P3 {
                                                                 P3 o(0, 0, 1);
P3F u = housev(q, o, q.z > 0 ? 1 : -1);
 lld x, y, z;
 P3 operator^(const P3 &b) const {
                                                                 auto pf = toP3F(p);
   return {y*b.z-b.y*z, z*b.x-b.z*x, x*b.y-b.x*y};
                                                                 auto np = pf - 2 * u * dot(u, pf) / dot(u, u);
                                                                 return P(np[0], np[1]);
  //Azimuthal angle (longitude) to x-axis. \in [-pi, pi]
                                                                } // project p onto the plane q^Tx = 0
 llf phi() const { return atan2(y, x); }
                                                                6.26 3D Skew Line Nearest Point
  //Zenith angle (latitude) to the z-axis. \in [0, pi]
 llf theta() const { return atan2(sqrt(x*x+y*y),z); }
                                                                • L_1: \mathbf{v}_1 = \mathbf{p}_1 + t_1 \mathbf{d}_1, L_2: \mathbf{v}_2 = \mathbf{p}_2 + t_2 \mathbf{d}_2
                                                                \cdot n = d_1 \times d_2
                                                                \cdot n_1 = d_1 \times n, n_2 = d_2 \times n
 P3 ver(P3 a, P3 b, P3 c) { return (b - a) ^ (c - a); }
                                                                m{\cdot} \ \ m{c}_1 = m{p}_1 + rac{(m{p}_2 - m{p}_1) \cdot m{n}_2}{m{d}_1 \cdot m{n}_2} m{d}_1, m{c}_2 = m{p}_2 + rac{(m{p}_1 - m{p}_2) \cdot m{n}_1}{m{d}_2 \cdot m{n}_1} m{d}_2
 lld volume(P3 a, P3 b, P3 c, P3 d) {
 return dot(ver(a, b, c), d - a);
                                                                       Delaunay [3a4ff1] - 1aee24/19ec42
                                                                   please ensure input points are unique *,
 P3 rotate_around(P3 p, llf angle, P3 axis) {
                                                                /* A triangulation such that no points will strictly
 llf s = sin(angle), c = cos(angle);
                                                                inside circumcircle of any triangle. C should be big
 P3 u = normalize(axis);
                                                                enough s.t. the initial triangle contains all points */
 return u*dot(u, p)*(1-c) + p * c + cross(u, p)*s;
                                                                #define L(i) ((i)==0 ? 2 : (i)-1)
                                                                #define R(i) ((i)==2 ? 0 : (i)+1)
 6.24 3D Convex Hull [01652a]
                                                                #define F3 for (int i = 0; i < 3; i++)
struct Face {
                                                                bool is_inf(P z) { return RE(z) \leftarrow -C || RE(z) \rightarrow= C; }
```

```
bool in_cc(const array<P,3> &p, P q) {
                                                                   P m = p[i] + p[j], d = rot90(p[j] - p[i]);
 i128 inf_det = 0, det = 0, inf_N, N;
                                                                   assert (norm(d) != 0);
                                                                   ls.emplace_back(m, m + d); // doubled coordinate
 F3 {
  if (is_inf(p[i]) && is_inf(q)) continue;
                                                                  } // HPI(ls)
  else if (is_inf(p[i])) inf_N = 1, N = -norm(q);
  else if (is_inf(q)) inf_N = -1, N = norm(p[i]);
                                                                6.29 kd Tree (Nearest Point)* [f733e5]
  else inf_N = 0, N = norm(p[i]) - norm(q);
  lld D = cross(p[R(i)] - q, p[L(i)] - q);
                                                                struct KDTree {
  inf_det += inf_N * D; det += N * D;
                                                                 struct Node {
                                                                  int x, y, x1, y1, x2, y2, id, f; Node *L, *R;
 return inf_det != 0 ? inf_det > 0 : det > 0;
                                                                 } tree[maxn], *root;
                                                                 lld dis2(int x1, int y1, int x2, int y2) {
P v[maxn];
                                                                  lld dx = x1 - x2, dy = y1 - y2;
struct Tri;
                                                                  return dx * dx + dy * dy;
struct E {
 Tri *t; int side;
                                                                 static bool cmpx(Node& a, Node& b) { return a.x<b.x; }</pre>
 E(Tri *t_=0, int side_=0) : t(t_), side(side_) {}
                                                                 static bool cmpy(Node& a, Node& b) { return a.y<b.y; }</pre>
                                                                 void init(vector<pair<int,int>> &ip) {
struct Tri {
                                                                  for (int i = 0; i < ssize(ip); i++)</pre>
 array<int,3> p; array<Tri*,3> ch; array<E,3> e;
                                                                   tie(tree[i].x, tree[i].y) = ip[i], tree[i].id = i;
 Tri(int a=0, int b=0, int c=0) : p{a, b, c}, ch{} {}
                                                                  root = build(0, (int)ip.size()-1, 0);
 bool has_chd() const { return ch[0] != nullptr; }
 bool contains(int q) const {
                                                                 Node* build(int L, int R, int d) {
 F3 if (ori(v[p[i]], v[p[R(i)]], v[q]) < 0)
                                                                  if (L>R) return nullptr;
   return false;
                                                                  int M = (L+R)/2;
  return true;
                                                                  nth_element(tree+L,tree+M,tree+R+1,d%2?cmpy:cmpx);
                                                                  Node &o = tree[M]; o.f = d \% 2;
                                                                  o.x1 = o.x2 = o.x; o.y1 = o.y2 = o.y;
 bool check(int q) const {
 return in_cc({v[p[0]], v[p[1]], v[p[2]]}, v[q]); }
                                                                  o.L = build(L, M-1, d+1); o.R = build(M+1, R, d+1);
                                                                  for (Node *s: {o.L, o.R}) if (s) {
} pool[maxn * 10], *it, *root;
                                                                   o.x1 = min(o.x1, s->x1); o.x2 = max(o.x2, s->x2);
/* SPLIT_HASH_HERE */
void link(const E &a, const E &b) {
                                                                   o.y1 = min(o.y1, s->y1); o.y2 = max(o.y2, s->y2);
 if (a.t) a.t->e[a.side] = b;
 if (b.t) b.t->e[b.side] = a;
                                                                  return tree+M;
                                                                 bool touch(int x, int y, lld d2, Node *r){
void flip(Tri *A, int a) {
 auto [B, b] = A->e[a]; /* flip edge between A,B */
                                                                  lld d = (lld) \operatorname{sqrt}(d2) + 1;
 if (!B || !A->check(B->p[b])) return;
                                                                  return x >= r->x1 - d && x <= r->x2 + d &&
         new (it++) Tri(A->p[R(a)], B->p[b], A->p[a]);
                                                                         y >= r->y1 - d \&\& y <= r->y2 + d;
 Tri *Y = new (it++) Tri(B->p[R(b)], A->p[a], B->p[b]);
 link(E(X, 0), E(Y, 0));
                                                                 using P = pair<lld, int>;
link(E(X, 1), A->e[L(a)]); link(E(X, 2), B->e[R(b)]);
link(E(Y, 1), B->e[L(b)]); link(E(Y, 2), A->e[R(a)]);
                                                                 void dfs(int x, int y, P &mn, Node *r) {
                                                                  if (!r || !touch(x, y, mn.first, r)) return;
A->ch = B->ch = {X, Y, nullptr};
flip(X, 1); flip(X, 2); flip(Y, 1); flip(Y, 2);
                                                                  mn = min(mn, P(dis2(r\rightarrow x, r\rightarrow y, x, y), r\rightarrow id));
                                                                  if (r->f == 1 ? y < r->y : x < r->x)
                                                                   dfs(x, y, mn, r\rightarrow L), dfs(x, y, mn, r\rightarrow R);
void add_point(int p) {
                                                                  else
Tri *r = root;
                                                                   dfs(x, y, mn, r\rightarrow R), dfs(x, y, mn, r\rightarrow L);
 while (r->has_chd()) for (Tri *c: r->ch)
 if (c && c->contains(p)) { r = c; break; }
                                                                 int query(int x, int y) {
 array<Tri*, 3> t; /* split into 3 triangles */
F3 t[i] = new (it++) Tri(r->p[i], r->p[R(i)], p);
                                                                  P mn(INF, -1); dfs(x, y, mn, root);
                                                                  return mn.second;
 F3 link(E(t[i], 0), E(t[R(i)], 1));
 F3 link(E(t[i], 2), r->e[L(i)]);
                                                                } tree;
6.30 kd Closest Pair (3D ver.)* [84d9eb]
 r->ch = t;
F3 flip(t[i], 2);
                                                                llf solve(vector<P> v) {
                                                                 shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,</pre>
auto build(const vector<P> &p) {
 it = pool; int n = (int)p.size();
                                                                  unordered_map<lld, int>>> m;
                                                                 llf d = dis(v[0], v[1]);
 vector<int> ord(n); iota(all(ord), 0);
shuffle(all(ord), mt19937(114514));
root = new (it++) Tri(n, n + 1, n + 2);
                                                                 auto Idx = [\&d] (llf x) \rightarrow lld {
                                                                  return round(x \star 2 / d) + 0.1; };
 copy_n(p.data(), n, v); v[n++] = P(-C, -C);
                                                                 auto rebuild_m = [&m, &v, &Idx](int k) {
 v[n++] = P(C * 2, -C); v[n++] = P(-C, C * 2);

for (int i : ord) add_point(i);
                                                                  m.clear();
                                                                  for (int i = 0; i < k; ++i)</pre>
 vector<array<int, 3>> res;
                                                                   m[Idx(v[i].x)][Idx(v[i].y)]
 for (Tri *now = pool; now != it; now++)
                                                                    [Idx(v[i].z)] = i;
  if (!now->has_chd()) res.push_back(now->p);
                                                                 }; rebuild_m(2);
                                                                 for (size_t i = 2; i < v.size(); ++i) {</pre>
 return res;
                                                                  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
       Build Voronoi [94f000]
                                                                     kz = Idx(v[i].z); bool found = false;
                                                                  for (int dx = -2; dx <= 2; ++dx) {
void build_voronoi_cells(auto &&p, auto &&res) {
 vector<vector<int>> adj(p.size());
                                                                   const lld nx = dx + kx;
 for (auto f: res) F3 {
                                                                   if (m.find(nx) == m.end()) continue;
  int a = f[i], b = f[R(i)];
                                                                   auto& mm = m[nx];
  if (a >= p.size() || b >= p.size()) continue;
                                                                   for (int dy = -2; dy <= 2; ++dy) {
                                                                    const lld ny = dy + ky;
if (mm.find(ny) == mm.end()) continue;
  adj[a].emplace_back(b);
 // use `adj` and `p` and HPI to build cells
                                                                    auto& mmm = mm[ny];
 for (size_t i = 0; i < p.size(); i++) {
  vector<Line> ls = frame; // the frame
                                                                    for (int dz = -2; dz <= 2; ++dz) {
                                                                      const lld nz = dz + kz;
  for (int j : adj[i]) {
                                                                      if (mmm.find(nz) == mmm.end()) continue;
```

```
const int p = mmm[nz];
                                                              induce(); vector<int> ns(lms.size());
     if (dis(v[p], v[i]) < d) {</pre>
                                                              for (int j = -1, nz = 0; int i : sa | is_lms) {
      d = dis(v[p], v[i]);
                                                               if (j >= 0) {
                                                                int len = min({n - i, n - j, lms[q[i] + 1] - i});
      found = true;
                                                                ns[q[i]] = nz += lexicographical_compare(
     }
    }
                                                                  begin(s) + j, begin(s) + j + len,
                                                                   begin(s) + i, begin(s) + i + len);
  }
  if (found) rebuild_m(i + 1);
                                                               j = i;
 else m[kx][ky][kz] = i;
                                                              }
                                                              ranges::fill(sa, 0); auto nsa = sais(ns);
                                                              for (auto x = c; int y : nsa | views::reverse)
return d;
                                                               y = lms[y], sa[--x[s[y]]] = y;
6.31 Simulated Annealing* [4e0fe5]
                                                              return induce(), sa;
llf anneal() {
                                                             // SPLIT_HASH_HERE sa[i]: sa[i]-th suffix is the
mt19937 rnd_engine(seed);
uniform_real_distribution<llf> rnd(0, 1);
                                                             // i-th lexicographically smallest suffix.
                                                             // hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
const llf dT = 0.001;
                                                             struct Suffix {
 // Argument p
                                                              int n; vector<int> sa, hi, rev;
llf S_cur = calc(p), S_best = S_cur;
                                                              Suffix(const auto &s) : n(int(s.size())),
 for (llf T = 2000; T > EPS; T -= dT) {
 // Modify p to p_prime
                                                               hi(n), rev(n) {
                                                               vector<int> _s(n + 1); // _s[n] = 0;
copy(all(s), begin(_s)); // s shouldn't contain 0
 const llf S_prime = calc(p_prime);
 const llf delta_c = S_prime - S_cur;
                                                               sa = sais(_s); sa.erase(sa.begin());
 llf prob = min((llf)1, exp(-delta_c / T));
                                                               for (int i = 0; i < n; ++i) rev[sa[i]] = i;
for (int i = 0, h = 0; i < n; ++i) {</pre>
 if (rnd(rnd_engine) <= prob)</pre>
   S_cur = S_prime, p = p_prime;
                                                                if (S_prime < S_best) // find min</pre>
  S_best = S_prime, p_best = p_prime;
                                                                hi[rev[i]] = h ? h-- : 0;
return S_best;
6.32 Triangle Centers* [adb146]
0 = ... // see min circle cover
                                                                  Suffix Array Tools* [8e08c8]
G = (A + B + C) / 3;
                                                             template <int LG = 20> struct SparseTableSA : Suffix {
H = G * 3 - 0 * 2; // orthogonal center
                                                              array<vector<int>, LG> mn;
llf a = abs(B - C), b = abs(A - C), c = abs(A - B);
                                                              SparseTableSA(const auto &s) : Suffix(s), mn{hi} {
I = (a * A + b * B + c * C) / (a + b + c);
                                                               for (int l = 0; l + 1 < LG; l++) { mn[l+1].resize(n);
for (int i = 0, len = 1 << l; i + len < n; i++)</pre>
// FermatPoint: minimizes sum of distance
// if max. angle >= 120 deg then vertex
                                                                 mn[l + 1][i] = min(mn[l][i], mn[l][i + len]);
// otherwise, make eq. triangle AB'C, CA'B, BC'A
// line AA', BB', CC' intersects at P
     Stringology
                                                              int lcp(int a, int b) {
7.1 Hash [ce7fad]
                                                               if (a == b) return n - a;
template <int P = 127, int Q = 1051762951>
                                                               a = rev[a] + 1, b = rev[b] + 1;
class Hash {
                                                               if (a > b) swap(a, b);
vector<int> h, p;
                                                               const int lg = __lg(b - a);
                                                                return min(mn[lg][a], mn[lg][b - (1 << lg)]);</pre>
Hash(const auto &s) : h(s.size()+1), p(s.size()+1) {
                                                              } // equivalent to lca on the kruskal tree
  for (size_t i = 0; i < s.size(); ++i)</pre>
                                                              pair<int,int> get_range(int x, int len) { // WIP
  h[i + 1] = add(mul(h[i], P), s[i]);
                                                               int a = rev[x] + 1, b = rev[x] + 1;
  generate(all(p), [x = 1, y = 1, this]() mutable {
                                                               for (int l = LG - 1; l >= 0; l--) {
   return y = x, x = mul(x, P), y; });
                                                                const int s = 1 << l;
                                                                 if (a + s <= n && mn[l][a] >= len) a += s;
int query(int l, int r) const { // 0-base [l, r)
                                                                if (b - s >= 0 && mn[l][b - s] >= len) b -= s;
 return sub(h[r], mul(h[l], p[r - l]));
                                                               return {b - 1, a};
                                                              } // if offline, solve get_range with DSU
      Suffix Array [ald8fe] - 9603d1/eb7a2f
                                                             7.4 Ex SAM* [58374b]
auto sais(const auto &s) {
const int n = (int)s.size(), z = ranges::max(s) + 1;
                                                             struct exSAM {
                                                              int len[maxn * 2], link[maxn * 2]; // maxlen, suflink
if (n == 1) return vector{0};
                                                              int next[maxn * 2][maxc], tot; // [0, tot), root = 0
vector<int> c(z); for (int x : s) ++c[x];
                                                              int ord[maxn * 2]; // topo. order (sort by length)
int cnt[maxn * 2]; // occurence
partial_sum(all(c), begin(c));
vector<int> sa(n); auto I = views::iota(0, n);
vector<bool> t(n); t[n - 1] = true;
                                                              int newnode() {
for (int i = n - 2; i >= 0; --i)
t[i] = (s[i]==s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);</pre>
                                                               fill_n(next[tot], maxc, 0);
return len[tot] = cnt[tot] = link[tot] = 0, tot++;
auto is_lms = views::filter([&t](int x) {
 return x && t[x] && !t[x - 1]; });
                                                              void init() { tot = 0, newnode(), link[0] = -1; }
 auto induce = [&] {
                                                              int insertSAM(int last, int c) {
 for (auto x = c; int y : sa)
                                                               int cur = next[last][c];
  if (y--) if (!t[y]) sa[x[s[y] - 1]++] = y;
                                                               len[cur] = len[last] + 1;
                                                               int p = link[last];
  for (auto x = c; int y : sa | views::reverse)
   if (y--) if (t[y]) sa[--x[s[y]]] = y;
                                                               while (p != -1 && !next[p][c])
                                                               next[p][c] = cur, p = link[p];
if (p == -1) return link[cur] = 0, cur;
};
vector<int> lms, q(n); lms.reserve(n);
                                                               int q = next[p][c];
for (auto x = c; int i : I | is_lms) {
 q[i] = int(lms.size());
                                                               if (len[p] + 1 == len[q]) return link[cur] = q, cur;
  lms.push_back(sa[--x[s[i]]] = i);
                                                               int clone = newnode();
                                                               for (int i = 0; i < maxc; ++i)</pre>
```

```
next[clone][i] = len[next[q][i]] ? next[q][i] : 0;
  len[clone] = len[p] + 1;
                                                             7.8
                                                                   Lyndon Factorization [d22cc9]
  while (p != -1 && next[p][c] == q)
                                                             // partition s = w[0] + w[1] + ... + w[k-1],
   next[p][c] = clone, p = link[p];
                                                             // w[0] >= w[1] >= ... >= w[k-1]
  link[link[cur] = clone] = link[q];
                                                             // each w[i] strictly smaller than all its suffix
  link[q] = clone;
                                                            void duval(const auto &s, auto &&report) {
  return cur;
                                                              for (int n = (int)s.size(), i = 0, j, k; i < n; ) {</pre>
                                                               for (j = i + 1, k = i; j < n \&\& s[k] <= s[j]; j++)
 void insert(const string &s) {
                                                               k = (s[k] < s[j] ? i : k + 1);
  int cur = 0;
                                                               // if (i < n / 2 && j >= n / 2) {
  for (char ch : s) {
                                                               // for min cyclic shift, call duval(s + s)
   int &nxt = next[cur][int(ch - 'a')];
                                                              // then here s.substr(i, n / 2) is min cyclic shift
   if (!nxt) nxt = newnode();
                                                               117
   cnt[cur = nxt] += 1;
                                                               for (; i <= k; i += j - k)</pre>
 }
                                                                report(i, j - k); // s.substr(l, len)
void build() {
                                                             }
                                                               // tested @ luogu 6114, 1368 & UVA 719
 queue<int> q; q.push(0);
                                                             7.9 Main Lorentz* [615b8f]
 while (!q.empty()) {
                                                             vector<pair<int, int>> rep[kN]; // 0-base [l, r]
   int cur = q.front(); q.pop();
                                                             void main_lorentz(const string &s, int sft = 0) {
  for (int i = 0; i < maxc; ++i)</pre>
                                                              const int n = s.size();
    if (next[cur][i]) q.push(insertSAM(cur, i));
                                                              if (n == 1) return;
                                                              const int nu = n / 2, nv = n - nu;
 vector<int> lc(tot);
                                                              const string u = s.substr(0, nu), v = s.substr(nu),
  for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
                                                                 ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend());
 partial_sum(all(lc), lc.begin());
                                                              main_lorentz(u, sft), main_lorentz(v, sft + nu);
  for (int i = 1; i < tot; ++i) ord[--lc[len[i]]] = i;</pre>
                                                              void solve() {
                                                              auto get_z = [](const vector<int> &z, int i) {
  for (int i = tot - 2; i >= 0; --i)
                                                               return (0 <= i and i < (int)z.size()) ? z[i] : 0; };</pre>
   cnt[link[ord[i]]] += cnt[ord[i]];
                                                              auto add_rep = [&](bool left, int c, int l, int k1,
                                                                 int k2) {
};
                                                               const int L = max(1, l - k2), R = min(l - left, k1);
7.5
    KMP [3727f3]
                                                               if (L > R) return;
vector<int> kmp(const auto &s) {
                                                               if (left) rep[l].emplace_back(sft + c - R, sft + c -
vector<int> f(s.size());
                                                                 L);
for (int i = 1, k = 0; i < (int)s.size(); ++i) {</pre>
                                                               else rep[l].emplace_back(sft + c - R - l + 1, sft + c
 while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
                                                                  - L - l + 1);
  f[i] = (k += (s[i] == s[k]));
                                                              for (int cntr = 0; cntr < n; cntr++) {</pre>
return f;
                                                               int l, k1, k2;
                                                               if (cntr < nu) {</pre>
vector<int> search(const auto &s, const auto &t) {
                                                                l = nu - cntr;
// return 0-indexed occurrence of t in s
                                                                k1 = get_z(z1, nu - cntr);
vector<int> f = kmp(t), r;
                                                                k2 = get_z(z2, nv + 1 + cntr);
for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
                                                               } else {
 while (k > 0 \&\& s[i] != t[k]) k = f[k - 1];
                                                                l = cntr - nu + 1;
 k += (s[i] == t[k]);
                                                                k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
  if (k == (int)t.size())
                                                                k2 = get_z(z4, (cntr - nu) + 1);
   r.push_back(i - t.size() + 1), k = f[k - 1];
                                                               if (k1 + k2 >= 1)
return r;
                                                                add_rep(cntr < nu, cntr, l, k1, k2);</pre>
}
7.6 Z value [6a7fd0]
vector<int> Zalgo(const string &s) {
                                                             7.10 BWT* [a8287e]
                                                             void BWT(char *ori, char *res) {
  // make ori -> ori + ori then build suffix array
vector<int> z(s.size(), s.size());
for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
 int j = clamp(r - i, 0, z[i - l]);
 for (; i + j < z[0] and s[i + j] == s[j]; ++j);
if (i + (z[i] = j) > r) r = i + z[l = i];
                                                             void iBWT(char *ori, char *res) {
                                                              vector<int> v[SIGMA], a;
                                                              const int len = strlen(ori); res[len] = 0;
                                                              for (int i = 0; i < len; i++) v[ori[i] - 'a'].pb(i);
for (int i = 0, ptr = 0; i < SIGMA; i++)</pre>
return z;
                                                             for (int j : v[i]) a.pb(j), ori[ptr++] = 'a' + i;
for (int i = 0, ptr = 0; i < len; i++)</pre>
     Manacher [c938a9]
vector<int> manacher(const string &S) {
                                                               res[i] = ori[a[ptr]], ptr = a[ptr];
const int n = (int)S.size(), m = n * 2 + 1;
vector<int> z(m);
string t = "."; for (char c: S) t += c, t += '.';
                                                                   Palindromic Tree* [0673ee]
for (int i = 1, l = 0, r = 0; i < m; ++i) {</pre>
                                                             struct PalindromicTree {
 z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
                                                              struct node {
 while (i - z[i] >= 0 \&\& i + z[i] < m) {
                                                               int nxt[26], f, len; // num = depth of fail link
   if (t[i - z[i]] == t[i + z[i]]) ++z[i];
                                                               int cnt, num;
                                                                                // = #pal_suffix of this node
   else break;
                                                               node(int l = 0) : nxt{}, f(0), len(l), cnt(0), num(0) {}
  if (i + z[i] > r) r = i + z[i], l = i;
                                                              vector<node> st; vector<char> s; int last, n;
}
                                                              void init() {
                                                               st.clear(); s.clear(); last = 1; n = 0;
return z; // the palindrome lengths are z[i] - 1
                                                               st.push_back(0); st.push_back(-1);
/* for (int i = 1; i + 1 < m; ++i) {
                                                               st[0].f = 1; s.push_back(-1);
  int l = (i - z[i] + 2) / 2, r = (i + z[i]) / 2;
  if (l != r) // [l, r) is maximal palindrome
```

int getFail(int x) {

```
while (s[n - st[x].len - 1] != s[n]) x = st[x].f;
 return x;
void add(int c) {
 s.push_back(c -= 'a'); ++n;
  int cur = getFail(last);
  if (!st[cur].nxt[c]) {
  int now = st.size();
   st.push_back(st[cur].len + 2);
  st[now].f = st[getFail(st[cur].f)].nxt[c];
  st[cur].nxt[c] = now;
  st[now].num = st[st[now].f].num + 1;
 last = st[cur].nxt[c]; ++st[last].cnt;
}
void dpcnt() { // cnt = #occurence in whole str
  for (int i = st.size() - 1; i >= 0; i--)
  st[st[i].f].cnt += st[i].cnt;
int size() { return st.size() - 2; }
} pt; /* string s; cin >> s; pt.init();
for (int i = 0; i < SZ(s); i++) {
int prvsz = pt.size(); pt.add(s[i]);
if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
 // pal @ [l,r]: s.substr(l, r-l+1)
8
    Misc
```

Spherical Coordinate $x = r \sin \theta \cos \phi$ $y = r \sin \theta \sin \phi$

Theorems

8.1



$$r = \sqrt{x^2 + y^2 + z^2}$$

$$\theta = \operatorname{acos}(z/\sqrt{x^2 + y^2 + z^2})$$

$$\phi = \operatorname{atan2}(y, x)$$

Spherical Cap

 $z = r \cos \theta$

- A portion of a sphere cut off by a plane.
- r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ :
- $\begin{aligned} &\text{Area} &= \pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta) (1 \cos \theta)^2/3. \\ &\text{Area} &= 2\pi r h = \pi (a^2 + h^2) = 2\pi r^2 (1 \cos \theta). \end{aligned}$

Sherman-Morrison formula

$$(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1 + v^{\mathsf{T}}A^{-1}u}$$

Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} =$ d(i), $L_{ij} = -c$ where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is $det(\tilde{L}_{11})$.
- The number of directed spanning tree rooted at r in G is $\det(\tilde{L}_{rr})$.

BEST Theorem

#{Eulerian circuits} = #{arborescences rooted at 1} $\cdot \prod_{v \in V} (\deg(v) - 1)!$

Random Walk on Graph

Let P be the transition matrix of a strongly connected directed graph, $\sum_{i} P_{i,j} = 1$. Let $F_{i,j}$ be the expected time to reach j from i. Let g_i be the expected time from i to i, G = diag(g) and J be a matrix all of 1, i.e. $J_{i,j} = 1$. Then, F = J - G + PF

First solve G: let $\pi P = \pi$ be a stationary distribution. Then $\pi_i g_i = 1$. The rank of I-P is n-1, so we first solve a special solution X such that (I-P)X=J-G and adjust X to F by $F_{i,j}=X_{i,j}-X_{j,j}$.

Tutte Matrix

For i < j, $d_{ij} = x_{ij}$ (in practice, a random number) if $(i,j) \in \mathit{E}$, otherwise $d_{ij}=0.$ For $i\geq j, d_{ij}=-d_{ji}.$ $rac{{\sf rank}(D)}{2}$ is the maximum matching.

Cayley's Formula

- Given a degree sequence d_1, d_2, \ldots, d_n for each labeled vertices, there're $rac{(n-2)!}{-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1, 2, \ldots, k$ belong to different components. Then $T_{n,k} = 0$

Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and $\sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k)$ holds for all $1 \le k \le n$

Havel-Hakimi algorithm

Find the vertex who has greatest degree unused, connect it with other great-

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and b_1, \ldots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i, k)$ holds for every $1 \le k \le n$

Fulkerson-Chen-Anstee theorem

A sequence $(a_1,b_1),\ldots,(a_n,b_n)$ of nonnegative integer pairs with $a_1\geq$ $\cdots \geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq a_i$ $\sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k)$ holds for every $1 \leq k \leq n$. **Euler's planar graph formula** V-E+F=C+1. $E \leq 3V-6$ (when $V \geq 3$)

Pick's theorem

For simple polygon, when points are all integer, we have A #{lattice points in the interior} + $\frac{1}{2}$ #{lattice points on the boundary} - 1

Matroid

 $\begin{array}{l} \bullet \ B \subseteq A \land A \in \mathcal{I} \Rightarrow B \in \mathcal{I}. \\ \bullet \ \text{If } A, B \in \mathcal{I} \text{ and } |A| > |B|, \text{then } \exists x \in A \setminus B, B \cup \{x\} \in \mathcal{I}. \end{array}$

Linear matroid $A \in I$ iff linear indep. I= forests of undirected graph Graphic matroid Colorful matroid (EX) Each color c has an upper bound R_c . Transversal matroid $A \in I$ iff \exists matching M whose right part is A. $\begin{array}{l} A\in I \text{ iff } G \text{ is connected after removing edges } A. \\ A\in I^* \text{ iff there is a basis} \subseteq E\setminus A \end{array}$ Bond matroid Dual matroid $A \in I' \text{ iff } A \in I \land |A| \le k$

Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2)$, find maximum $S\in I_1\cap I_2$. For each iteration, build the directed graph and find a shortest path from s to t.

• $s \rightarrow x : S \sqcup \{x\} \in I_1$ $x \to t : S \sqcup \{x\} \in I_2$

Truncated matroid

- $y \to x : S \setminus \{y\} \sqcup \{x\} \in I_1$ (y is in the unique circuit of $S \sqcup \{x\}$) $x \to y : S \setminus \{y\} \sqcup \{x\} \in I_2$ (y is in the unique circuit of $S \sqcup \{x\}$)

Alternate the path, and |S| will increase by 1. In each iteration, |E|=O(RN), where $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)), N=|G|$. For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Find the shortest path by Bellman-Ford. The maximum iteration of Bellman-Ford is 2R+1.

Dual of LP

2 44. 0. 2.	
Primal	Dual
Maximize $c^{T}x$ s.t. $Ax \leq b, x \geq 0$	Minimize $b^{T}y$ s.t. $A^{T}y \geq c$, $y \geq 0$
Maximize $c^{T}x$ s.t. $Ax \leq b$	Minimize $b^{T}y$ s.t. $A^{T}y = c, y \geq 0$
Maximize $c^{T}x$ s.t. $Ax = b, x > 0$	Minimize $b^{T} u$ s.t. $A^{T} u \geq c$

Dual of Min Cost b-Flow

- Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .
- If all w_{uv} are integers, then optimal solution can happen when all p_u are

$$\min \sum_{uv} w_{uv} f_{uv} \text{ s.t. } -f_{uv} \geq -c_{uv}, \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u}$$

$$\Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \text{ s.t. } p_{u} \geq 0$$

Minimax Theorem

Let $f:X\times Y\to\mathbb{R}$ be continuous where $X\subseteq\mathbb{R}^n,Y\subseteq\mathbb{R}^m$ are compact and convex. If $f(\cdot,y):X\to\mathbb{R}$ is concave for fixed y, and $f(x,\cdot):Y\to\mathbb{R}$ is convex for fixed x, then $\max_{x \in X} \min_{y \in Y} f(x,y) = \min_{y \in Y} \max_{x \in X} f(x,y)$, e.g. $f(x,y) = x^{\mathsf{T}} Ay$ for

zero-sum matrix game. Parallel Axis Theorem

The second moment of area is $I_z=\iint x^2+y^2\mathrm{d}A.$ $I_{z'}=I_z+Ad^2$ where d is the distance between two parallel axis z,z'.

Stable Marriage

```
1: Initialize m \in M and w \in W to free
    while \exists free man m who has a woman w to propose to do
        w \leftarrow first woman on m's list to whom m has not yet proposed
        if \exists some pair (m', w) then
5:
           if w prefers m to m^\prime then
6:
7:
               m' \leftarrow \textit{free}
               (m,w) \leftarrow \mathsf{engaged}
8:
           end if
        else
10:
            (m, w) \leftarrow \mathsf{engaged}
        end if
12: end while
```

Weight Matroid Intersection* [d00ee8]

```
struct Matroid {
 Matroid(bitset<N>); // init from an independent set
bool can_add(int); // check if break independence
 Matroid remove(int); // removing from the set
auto matroid_intersection(const vector<int> &w) {
 const int n = (int)w.size(); bitset<N> S;
 for (int sz = 1; sz <= n; sz++) {</pre>
  Matroid M1(S), M2(S); vector<vector<pii>>> e(n + 2);
  for (int j = 0; j < n; j++) if (!S[j]) {</pre>
   if (M1.can_add(j)) e[n].eb(j, -w[j]);
if (M2.can_add(j)) e[j].eb(n + 1, 0);
  for (int i = 0; i < n; i++) if (S[i]) {</pre>
   Matroid T1 = M1.remove(i), T2 = M2.remove(i);
   for (int j = 0; j < n; j++) if (!S[j]) {</pre>
    if (T1.can_add(j)) e[i].eb(j, -w[j]);
    if (T2.can_add(j)) e[j].eb(i, w[i]);
  \} // maybe implicit build graph for more speed
  vector<pii> d(n + 2, \{INF, 0\}); d[n] = \{0, 0\};
  vector<int> prv(n + 2, -1);
  // change to SPFA for more speed, if necessary
  for (int upd = 1; upd--; )
```

for (int u = 0; u < n + 2; u++)

```
for (auto [v, c] : e[u]) {
                                                                 Isect(x, erase(y));
     pii x(d[u].first + c, d[u].second + 1);
     if (x < d[v]) d[v] = x, prv[v] = u, upd = 1;
                                                               lld Query(lld x) { // default chmax
                                                                auto l = *lower_bound(x); // to chmin:
return l.a * x + l.b; // modify the 2 "<>"
  if (d[n + 1].first >= INF) break;
  for (int x = prv[n+1]; x!=n; x = prv[x]) S.flip(x);
                                                               }
  // S is the max-weighted independent set w/ size sz
                                                              8.8
                                                                    Min Plus Convolution [464dcd]
 return S;
                                                              // a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
} // from Nacl
                                                              vector<int> min_plus_convolution(auto &a, auto &b) {
8.4 Bitset LCS [4155ab]
                                                               const int n = (int)a.size(), m = (int)b.size();
cin >> n >> m;
                                                               vector<int> c(n + m - 1, numeric_limits<int>::max());
for (int i = 1, x; i \le n; ++i)
                                                               auto dc = [&](auto Y, int l, int r, int jl, int jr) {
cin >> x, p[x].set(i);
                                                                if (l > r) return;
for (int i = 1, x; i <= m; ++i) {
  cin >> x, (g = f) |= p[x];
                                                                int mid = (l + r) / 2, from = -1, &best = c[mid];
                                                                for (int j = jl; j <= jr; j++)</pre>
 f.shiftLeftByOne(), f.set(0);
                                                                 if (int i = mid - j; i >= 0 && i < n)</pre>
 ((f = g - f) ^= g) \&= g;
                                                                  if (best > a[i]+b[j]) best = a[i]+b[j], from = j;
                                                                Y(Y, l, mid-1, jl, from); Y(Y, mid+1, r, from, jr);
cout << f.count() << '\n';</pre>
8.5 Prefix Substring LCS [7d8faf]
                                                               return dc(dc, 0, n-1+m-1, 0, m-1), c;
void all_lcs(string S, string T) { // 0-base
                                                              8.9 SMAWK [f37761]
 vector<size_t> h(T.size()); iota(all(h), 1);
 for (size_t a = 0; a < S.size(); ++a) {</pre>
                                                             // For all 2x2 submatrix:
  for (size_t c = 0, v = 0; c < T.size(); ++c)</pre>
                                                              // If M[1][0] < M[1][1], M[0][0] < M[0][1]
   if (S[a] == T[c] || h[c] < v) swap(h[c], v);</pre>
                                                              // If M[1][0] == M[1][1], M[0][0] <= M[0][1]
                                                              // M[i][ans_i] is the best value in the i-th row
  // here, LCS(s[0, a], t[b, c]) =
  // c - b + 1 - sum([h[i] > b] | i <= c)
                                                             VI smawk(int N, int M, auto &&select) {
}
                                                               auto dc = [&](auto self, const VI &r, const VI &c) {
} // test @ yosupo judge
                                                                if (r.empty()) return VI{};
                                                                const int n = (int)r.size(); VI ans(n), nr, nc;
8.6 Convex 1D/1D DP* [938911]
                                                                for (int i : c) {
struct S { int i, l, r; };
auto solve(int n, int k, auto &w) {
                                                                 while (!nc.empty() &&
                                                                   select(r[nc.size() - 1], nc.back(), i))
 vector < int64_t > dp(n + 1); dp[0] = 0;
                                                                  nc.pop_back();
 auto f = [&](int l, int r) -> int64_t {
  if (r - l > k) return -INF;
                                                                 if (int(nc.size()) < n) nc.push_back(i);</pre>
  return dp[l] + w(l + 1, r);
                                                                for (int i = 1; i < n; i += 2) nr.push_back(r[i]);</pre>
                                                                const auto na = self(self, nr, nc);
 deque<S> dq; dq.emplace_back(0, 1, n);
                                                                for (int i = 1; i < n; i += 2) ans[i] = na[i >> 1];
 for (int i = 1; i <= n; ++i) {
                                                                for (int i = 0, j = 0; i < n; i += 2) {
  dp[i] = f(dq.front().i, i);
                                                                 ans[i] = nc[j];
  while (!dq.empty() && dq.front().r <= i)</pre>
                                                                 const int end = i + 1 == n ? nc.back() : ans[i + 1];
  dq.pop_front();
                                                                 while (nc[j] != end)
  dq.front().l = i + 1;
                                                                  if (select(r[i], ans[i], nc[++j])) ans[i] = nc[j];
  while (!dq.empty() &&
    f(i, dq.back().l) >= f(dq.back().i, dq.back().l))
                                                                return ans;
   dq.pop_back();
  int p = i + 1;
                                                               VI R(N), C(M); iota(all(R), 0), iota(all(C), 0);
  if (!dq.empty()) {
                                                               return dc(dc, R, C);
   auto [j, l, r] = dq.back();
   for (int s = 1 << 20; s; s >>= 1)
   if (l+s <= n && f(i, l+s) < f(j, l+s)) l += s;
dq.back().r = l; p = l + 1;</pre>
                                                              bool min_plus_conv_select(int r, int u, int v) {
                                                               auto f = [](int i, int j) {
                                                                if (0 <= i - j && i - j < n) return b[j] + a[i - j];</pre>
                                                                return 2100000000 + (i - j);
 if (p <= n) dq.emplace_back(i, p, n);</pre>
                                                              return f(r, u) > f(r, v);
} // if f(r, v) is better than f(r, u), return true
 return dp;
  // test @ tioj 烏龜疊疊樂
                                                              8.10 De-Bruijn [aa7700]
8.7 ConvexHull Optimization [b4318e]
                                                              vector<int> de_bruijn(int k, int n) {
  // return cyclic string of len k^n s.t. every string
struct L {
 mutable lld a, b, p;
                                                               // of len n using k char appears as a substring.
 bool operator<(const L &r) const {</pre>
                                                               vector<int> aux(n + 1), res;
  return a < r.a; /* here */ }
                                                               auto db = [&](auto self, int t, int p) -> void {
 bool operator<(lld x) const { return p < x; }</pre>
                                                                if (t <= n)
                                                                 for (int i = aux[t - p]; i < k; ++i, p = t)</pre>
lld Div(lld a, lld b) {
  return a / b - ((a ^ b) < 0 && a % b); }</pre>
                                                                  aux[t] = i, self(self, t + 1, p);
                                                                else if (n % p == 0) for (int i = 1; i <= p; ++i)
struct DynamicHull : multiset<L, less<>>> {
                                                                 res.push_back(aux[i]);
 static const lld kInf = 1e18;
                                                               };
 bool Isect(iterator x, iterator y) {
                                                               return db(db, 1, 1), res;
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a)
  x->p = x->b > y->b ? kInf : -kInf; /* here */
                                                              8.11 Josephus Problem [7f9ceb]
                                                              lld f(lld n, lld m, lld k) { // n 人每次隔 m-1 個殺 lld s = (m - 1) % (n - k); // O(k)
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
                                                               for (lld i = n - k + 1; i <= n; i++) s = (s + m) % i;
 void Insert(lld a, lld b) {
                                                               return s;
 auto z = insert({a, b, 0}), y = z++, x = y;
                                                              lld kth(lld n, lld m, i128 k) { // died at kth
  while (Isect(y, z)) z = erase(z);
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
                                                               if (m == 1) return k; // O(m log(n))
                                                               for (k = k*m+m-1; k >= n; k = k-n + (k-n)/(m-1));
  while ((y = x) != begin() && (--x)->p >= y->p)
```

```
return k;
} // k and result are 0-based, test @ CF 101955
8.12 N Queens Problem* [31f83e]
void solve(VI &ret, int n) { // no sol when n=2,3
if (n % 6 == 2) {
 for (int i = 2; i <= n; i += 2) ret.push_back(i);</pre>
  ret.push_back(3); ret.push_back(1);
 for (int i = 7; i <= n; i += 2) ret.push_back(i);</pre>
 ret.push_back(5);
} else if (n % 6 == 3) {
 for (int i = 4; i <= n; i += 2) ret.push_back(i);</pre>
 ret.push_back(2);
 for (int i = 5; i <= n; i += 2) ret.push_back(i);</pre>
 ret.push_back(1); ret.push_back(3);
} else {
 for (int i = 2; i <= n; i += 2) ret.push_back(i);</pre>
 for (int i = 1; i <= n; i += 2) ret.push_back(i);</pre>
8.13
       Manhattan MST [1008bc]
vector<array<int, 3>> manhattanMST(vector<P> ps) {
vector<int> id(ps.size()); iota(all(id), 0);
vector<array<int, 3>> edges;
for (int k = 0; k < 4; k++) {
 sort(all(id), [&](int i, int j) {
  return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y; });</pre>
 map<int, int> sweep;
 for (int i : id) {
  for (auto it = sweep.lower_bound(-ps[i].y);
     it != sweep.end(); sweep.erase(it++)) {
    if (P d = ps[i] - ps[it->second]; d.y > d.x) break;
    else edges.push_back({d.y + d.x, i, it->second});
  sweep[-ps[i].y] = i;
 for (P &p : ps)
  if (k \& 1) p.x = -p.x;
  else swap(p.x, p.y);
return edges; // [{w, i, j}, ...]
} // test @ yosupo judge
8.14 Binary Search On Fraction [ff3abd]
struct Q {
lld p, q; // p / q
Q go(Q b, lld d) { return {p + b.p*d, q + b.q*d}; }
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(lld N, auto &&pred) {
Q lo{0, 1}, hi{1, 0};
if (pred(lo)) return lo;
assert(pred(hi));
bool dir = 1, L = 1, H = 1;
 for (; L || H; dir = !dir) {
 lld len = 0, step = 1;
 for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
  if (Q mid = hi.go(lo, len + step);
    mid.p > N || mid.q > N || dir ^ pred(mid))
   t++;
  else len += step;
 swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
return dir ? hi : lo;
8.15
      Cartesian Tree [2ed09d]
auto CartesianTree(const auto &a) {
const int n = (int)a.size(); vector<int> pa(n+1, -1);
for (int i = 1; i < n; i++) {</pre>
 int &p = pa[i] = i - 1, l = n;
 while (p != -1 && a[i] < a[p])
  tie(l, pa[l], p, pa[p]) = tuple(p, p, pa[p], i);
return pa.pop_back(), pa;
} // root is minimum
8.16 Nim Product [4ac1ce]
#define rep(i, r) for (int i = 0; i < r; i++)
struct NimProd {
llu bit_prod[64][64]{}, prod[8][8][256][256]{};
NimProd() {
 rep(i, 64) rep(j, 64) if (i & j) {
```

```
int a = lowbit(i & j);
  bit_prod[i][j] = bit_prod[i ^ a][j] ^
  bit_prod[(i ^ a) | (a-1)][(j ^ a) | (i & (a-1))];
} else bit_prod[i][j] = 1ULL << (i | j);</pre>
rep(e, 8) rep(f, 8) rep(x, 256) rep(y, 256)
 rep(i, 8) if (x >> i & 1) rep(j, 8) if (y >> j & 1)
  prod[e][f][x][y] ^= bit_prod[e * 8 + i][f * 8 + j];
llu operator()(llu a, llu b) const {
llu r = 0;
rep(e, 8) rep(f, 8)
 r ^= prod[e][f][a >> (e*8) & 255][b >> (f*8) & 255];
```

8.17 Grid

