Contents

1	Basi	6													1
•	1.1	vimrc													
	1.2	Debug Macro													
	1.3	Increase Stack													
	1.4	Pragma Optimization													
	1.5	IO Optimization													 . 2
_															_
2		Structure													2
	2.1	Dark Magic													
	2.3	LiChao Segment Tree													
	2.4	Treap													
	2.5	Linear Basis													
	2.6	Binary Search On Segment													
3	Gra														3
	3.1	BCC Edge													
	3.2	BCC Vertex													
	3.3	2-SAT (SCC)													
	3.4 3.5	Lowbit Decomposition													
	3.6	MaxClique													
	3.7	Virtural Tree													
	3.8	Centroid Decomposition													
	3.9	Tree Hashing													
		Minimum Mean Cycle													
	3.11	Mo's Algorithm on Tree													
		Minimum Steiner Tree													
		Directed Minimum Spanning													
		Manhattan Minimum Spann													
		Dominator Tree													
	5.16	Edge Coloring			•		٠.	٠	 ٠	 •	•	 •	•	•	 . 9
4	Mat	ching & Flow													9
٠	4.1	Kuhn Munkres													-
	4.2	Bipartite Matching													
	4.3	General Graph Matching													
	4.4	Minimum Weight Matching	(Cliq	Ιυe	ve	rsi	on)								 . 10
	4.5	Minimum Cost Circulation .													 . 10
	4.6	Flow Models													
	4.7	Dinic													
	4.8	Minimum Cost Maximum Flo													
	4.9	GomoryHu Tree													
		Global Min-Cut													
	7.11	Dijkstra Cost Flow			•			•	 •	 •	•	 •	•	•	 12
5	Mat	h													13
5	Mat 5.1	h $\lfloor rac{n}{i} floor$ Enumeration $\ldots \ldots$													
5	5.1 5.2	$\lfloor \frac{n}{i} \rfloor$ Enumeration ax+by=gcd													 13 13
5	5.1 5.2 5.3	$\lfloor \frac{n}{i} \rfloor$ Enumeration ax+by=gcd Pollard Rho		 				:	 	 		 			 13 13 13
5	5.1 5.2 5.3 5.4	$ \begin{array}{c} \lfloor \frac{n}{i} \rfloor \text{ Enumeration } \dots \\ \text{ax+by=gcd } \dots \\ \text{Pollard Rho } \dots \\ \text{Pi Count (Linear Sieve) } \dots \end{array} $:	 	 		 			 13 13 13 13
5	5.1 5.2 5.3					 			 	 		 			 13 13 13 13 13
5	5.1 5.2 5.3 5.4	$ \begin{array}{cccc} \lfloor \frac{n}{t} \rfloor & \text{Enumeration} & \dots & \\ \text{ax+by=gcd} & \dots & \\ \text{Pollard Rho} & \dots & \\ \text{Pi Count (Linear Sieve)} & \dots & \\ \text{Strling Number} & \dots & \\ \text{5.5.1} & \text{First Kind} & \dots & \\ \end{array} $				 			 	 		 			 13 13 13 13 13 13
5	5.1 5.2 5.3 5.4 5.5	$ \begin{array}{cccc} \lfloor \frac{n}{T} \rfloor & \text{Enumeration} & \dots & \\ \text{ax+by=gcd} & \dots & \dots & \\ \text{Pollard Rho} & \dots & \dots & \\ \text{Pi Count (Linear Sieve)} & \dots & \\ \text{Strling Number} & \dots & \dots & \\ \text{5.5.1} & \text{First Kind} & \dots & \\ \text{5.5.2} & \text{Second Kind} & \dots & \dots & \\ \end{array} $				 			 	 		 			 13 13 13 13 13 13 13
5	5.1 5.2 5.3 5.4 5.5	$ \begin{array}{cccc} \lfloor \frac{n}{T} \rfloor & \text{Enumeration} & \dots & \\ \text{ax+by=gcd} & \dots & \dots & \\ \text{Pollard Rho} & \dots & \dots & \\ \text{Pi Count (Linear Sieve)} & \dots & \\ \text{Strling Number} & \dots & \dots & \\ \text{5.5.1} & \text{First Kind} & \dots & \\ \text{5.5.2} & \text{Second Kind} & \dots & \\ \text{Range Sieve} & \dots & \dots & \\ \end{array} $							 	 		 			 13 13 13 13 13 13 13 13
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7	$ \begin{array}{cccc} \lfloor \frac{n}{t} \rfloor & \text{Enumeration} & \dots & \\ \text{ax+by=gcd} & \dots & \dots & \\ \text{Pollard Rho} & \dots & \\ \text{Pi Count (Linear Sieve)} & \dots & \\ \text{Strling Number} & \dots & \dots & \\ \text{5.5.1} & \text{First Kind} & \dots & \dots & \\ \text{5.5.2} & \text{Second Kind} & \dots & \\ \text{Range Sieve} & \dots & \dots & \\ \text{Miller Rabin} & \dots & \dots & \dots & \\ \end{array} $							 	 		 			 13 13 13 13 13 13 13 13 13
5	5.1 5.2 5.3 5.4 5.5	$ \begin{array}{cccc} \lfloor \frac{n}{T} \rfloor & \text{Enumeration} & \dots & \\ \text{ax+by=gcd} & \dots & \dots & \\ \text{Pollard Rho} & \dots & \dots & \\ \text{Pi Count (Linear Sieve)} & \dots & \\ \text{Strling Number} & \dots & \dots & \\ \text{5.5.1} & \text{First Kind} & \dots & \\ \text{5.5.2} & \text{Second Kind} & \dots & \\ \text{Range Sieve} & \dots & \dots & \\ \end{array} $							 	 		 			 13 13 13 13 13 13 13 13 13
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9	$ \begin{array}{cccc} \lfloor \frac{n}{t} \rfloor & \text{Enumeration} & \dots & \\ \text{ax+by=gcd} & \dots & \dots & \\ \text{Pollard Rho} & \dots & \\ \text{Pi Count (Linear Sieve)} & \dots & \\ \text{Strling Number} & \dots & \dots & \\ \text{5.5.1 First Kind} & \dots & \dots & \\ \text{5.5.2 Second Kind} & \dots & \dots & \\ \text{Range Sieve} & \dots & \dots & \\ \text{Miller Rabin} & \dots & \dots & \\ \text{Extended Euler} & \dots & \dots & \\ \end{array} $							 	 		 			 13 13 13 13 13 13 13 13 13 14 14
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10	$ \begin{bmatrix} \frac{n}{T} \end{bmatrix} \text{Enumeration} \\ \text{ax+by=gcd} \\ \text{Pollard Rho} \\ \text{Pi Count (Linear Sieve)} \\ \text{Strling Number} \\ \text{5.5.1 First Kind} \\ \text{5.5.2 Second Kind} \\ \text{Range Sieve} \\ \text{Miller Rabin} \\ \text{Extended Euler} \\ \text{Gauss Elimination} $							 	 					13 13 13 13 13 13 13 13 14 14 14
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12	$ \begin{bmatrix} \frac{n}{T} \end{bmatrix} \text{Enumeration} \\ \text{ax+by=gcd} \\ \text{Pollard Rho} \\ \text{Pi Count (Linear Sieve)} \\ \text{Strling Number} \\ \text{S.5.1 First Kind} \\ \text{5.5.2 Second Kind} \\ \text{Range Sieve} \\ \text{Miller Rabin} \\ \text{Extended Euler} \\ \text{Gauss Elimination} \\ \text{Fast Fourier Transform} \\ \text{Chinese Remainder} \\ \text{Berlekamp Massey} \\ \dots$							 			 			13 13 13 13 13 13 13 13 14 14 14 14 14
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13	\begin{align*} \begin{align*} \begin{align*} \left alig										 			13 13 13 13 13 13 13 13 14 14 14 14 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14	\begin{align*} \begin{align*} \begin{align*} \left alig										 			13 13 13 13 13 13 13 13 14 14 14 14 15 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15	[n]/2 Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT										 			13 13 13 13 13 13 13 13 14 14 14 14 15 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16	\begin{align*} \begin{align*} \begin{align*} \left(\begi										 			13 13 13 13 13 13 13 13 13 14 14 14 14 15 15 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17	[**] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum													133 133 133 133 133 134 144 144 145 155 156 166 166 166
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18	[n] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum ExtendedFloorSum													133 133 133 134 144 144 145 155 166 166 166 17
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19	[n] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue													133 133 133 134 144 144 145 155 166 166 177 17
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20	[*] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn													133 133 133 133 133 134 144 144 145 155 166 166 177 177 17
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.14 5.15 5.16 5.17 5.18 5.19 5.10 5.11 5.12 5.13 5.14 5.15 5.10 5.11 5.12 5.13 5.14 5.15	[n] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Simplex													133 133 133 133 133 134 144 144 144 155 166 166 167 177 177 177 177 177 177
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.14 5.15 5.16 5.17 5.18 5.19 5.10 5.11 5.12 5.13 5.14 5.15 5.10 5.11 5.12 5.13 5.14 5.15	[*] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Construction													133 133 133 133 133 134 144 144 144 155 166 166 167 177 177 177 177 177 177
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.15 5.16 5.17 5.18 5.19 5.20 5.20 5.21 5.22 5.22	[n] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Simplex													133 133 133 133 133 134 144 144 144 155 166 166 177 177 177 18
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22 5.23 5.24	[*] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Charateristic Polynomial Partition Number													133 133 133 133 133 134 144 145 155 166 167 177 177 177 188 18
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22 5.23 5.24 Geo	[*] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Charateristic Polynomial Partition Number metry													133 133 133 133 134 144 144 145 155 156 166 167 177 177 177 178 188 18
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22 5.23 5.24 6.27 6.29 6.20 6.20 6.20 6.20 6.20 6.20 6.20 6.20	1/1 Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Charateristic Polynomial Partition Number													133 133 133 133 133 134 144 144 145 155 156 166 177 177 178 188 18
	5.1 5.2 5.3 5.4 5.5 5.10 5.11 5.12 5.13 5.14 5.15 5.19 5.20 5.21 5.23 5.24 Geo 6.1 6.2	1/1 Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Charateristic Polynomial Partition Number metry Basic Geometry 2D Convex Hull													133 133 133 133 133 134 144 144 144 155 156 167 177 177 177 177 177 178 188 18
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22 5.23 5.24 6.27 6.29 6.20 6.20 6.20 6.20 6.20 6.20 6.20 6.20	□ □													133 133 133 133 133 133 134 144 144 155 166 167 177 177 177 188 18 18 18 18 18
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.14 5.15 5.16 5.17 5.20 5.21 5.22 5.23 5.24 6.2 6.3	1/1 Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Charateristic Polynomial Partition Number metry Basic Geometry 2D Convex Hull													133 133 133 133 133 134 144 144 145 155 166 167 177 177 177 188 188 188 188 19
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.19 5.20 5.21 5.22 5.23 5.24 Geo 6.1 6.2 6.3 6.4	[*] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Charateristic Polynomial Partition Number metry Basic Geometry 2D Convex Hull 3D Convex Hull 2D Farthest Pair													133 133 133 133 133 134 144 144 144 145 155 156 166 177 177 188 188 188 188 189 199 199
	5.1 5.2 5.3 5.4 5.5 5.10 5.11 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 6.2 6.3 6.4 6.5 6.6 6.7	□ 1/2 Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Construction Charateristic Polynomial Partition Number metry Basic Geometry 2D Convex Hull 3D Convex Hull 2D Closest Pair kD Closest Pair kD Closest Pair (3D ver.) Simulated Annealing													133 133 133 133 133 133 134 144 144 155 156 166 166 167 177 177 177 177 178 188 188 189 199 199 199 199
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.14 5.15 5.16 5.17 5.19 5.20 5.21 5.22 5.24 6.3 6.4 6.6 6.6 6.7 6.8	1 / 1 / 2 / 2 Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Charateristic Polynomial Partition Number metry Basic Geometry 2D Convex Hull 2D Convex Hull 2D Farthest Pair 2D Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection													133 133 133 133 133 133 134 144 144 155 156 166 167 177 177 178 188 188 189 199 199 199 199 199 199
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.19 5.20 5.21 5.22 5.23 6.4 6.5 6.6 6.7 6.6 6.7 6.8 6.9	[*] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Charateristic Polynomial Partition Number metry Basic Geometry 2D Convex Hull 3D Convex Hull 2D Farthest Pair 4D Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum													133 133 133 133 133 134 144 144 145 155 166 167 177 177 177 177 177 179 199 199 199 19
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.22 5.21 5.22 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10	1/1 Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Construction Simplex Construction Simplex Construction Simplex Construction Simplex Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Intersection of line and Circ													133 133 133 133 133 133 134 144 144 144
	5.1 5.2 5.3 5.4 5.5 5.10 5.11 5.15 5.16 5.17 5.18 5.19 5.20 5.21 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11	1 / 1 / 2 / 2 Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Construction Simplex Charateristic Polynomial Partition Number Partition Number metry Basic Geometry 2D Convex Hull 3D Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Intersection of Polygon and													133 133 133 133 133 134 144 144 145 155 156 166 167 177 177 177 177 177 177 177 17
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.15 5.16 5.17 5.18 5.19 5.20 5.21 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 6.12 6.12 6.13 6.14 6.15 6.15 6.15 6.15 6.15 6.15 6.15 6.15	1/2 Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Charateristic Polynomial Partition Number metry Basic Geometry 2D Convex Hull 3D Convex Hull 3D Convex Hull 2D Closest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Intersection of Floygon and Intersection of Two Circle													133 133 133 133 133 133 134 144 144 155 156 166 166 167 177 177 177 177 177 177 17
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.14 5.15 5.16 5.17 5.19 5.20 5.21 5.22 5.24 6.6 6.7 6.8 6.9 6.10 6.11 6.12 6.13	1/2 Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Charateristic Polynomial Partition Number metry Basic Geometry 2D Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Intersection of Polygon and Intersection of Two Circle Tangent line of Two Circle													133 133 133 133 133 133 134 144 144 155 166 166 177 177 177 177 178 188 188 189 199 199 199 200 200 200 200
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.19 5.20 5.21 5.22 5.23 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 6.12 6.12 6.13 6.13 6.14 6.14 6.15 6.15 6.16 6.16 6.16 6.16 6.16 6.16	[*] Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Charateristic Polynomial Partition Number metry Basic Geometry 2D Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Intersection of Polygon and Intersection of Two Circle Tangent line of Two Circle Minimum Covering Circle <td></td> <td>133 133 133 133 133 134 144 144 145 155 166 166 177 177 177 177 177 177 179 199 199 199</td>													133 133 133 133 133 134 144 144 145 155 166 166 177 177 177 177 177 177 179 199 199 199
	5.1 5.2 5.3 5.4 5.5 5.10 5.12 5.13 5.14 5.15 5.15 5.17 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.10 6.11 6.12 6.14 6.15	1/2 Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.5.1 First Kind 5.5.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum ExtendedFloorSum Quadratic residue De-Bruijn Simplex Charateristic Polynomial Partition Number metry Basic Geometry 2D Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Intersection of Polygon and Intersection of Two Circle Tangent line of Two Circle													133 133 133 133 133 133 134 144 144 145 155 156 166 177 177 177 177 188 188 189 199 199 190 200 200 200 200 200 200 200 200 200 2

```
7 Stringology
7.1 Hash . .
 Suffix Array
Suffix Automaton
KMP
Z value
                   22
7.5
 7.6
 7.7
 8 Misc
8.1 Theorems . .
  8.1.7
   8.2.1 totally monotone (concave/convex) . . . . . . . . . . . . . . .
  8.2.2 monge condition (concave/convex) . . . . . . . . . . . . . . .
8.5 Josephus Problem
8.6 Cactus Matching
8.7 Tree Knapsack
8.8 N Queens Problem
1
 Basic
```

1.1 vimrc

```
se is nu bs=2 ru mouse=a encoding=utf-8 ls=2
se cin cino+=j1 et sw=4 sts=4 tgc sc hls
syn on
colorscheme desert
filetype indent on
inoremap \{ < CR > \{ < CR > \} < ESC > 0 \}
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -std=c++17 -
    DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
    Wconversion -fsanitize=address,undefined -g && echo
     success<CR>
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 &&
     echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
```

1.2 Debug Macro

```
#ifdef KISEKI
#define safe cerr<<__PRETTY_FUNCTION__\
<<" line "<<__LINE__<<" safe\n"</pre>
#define debug(a...) qwerty(#a, a)
#define orange(a...) dvorak(#a, a)
using std::cerr;
template <typename ...T>
void qwerty(const char *s, T ...a) {
  cerr << "\e[1;32m(" << s << ") = (";
  int cnt = sizeof...(T);</pre>
   (..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
}
template <typename Iter>
void dvorak(const char *s, Iter L, Iter R) {
  cerr << "\e[1;32m[ " << s << " ] = [ ";</pre>
  for (int f = 0; L != R; ++L)
    cerr << (f++? ", ": "") << *L;
    cerr << "]\e[0m\n";
}
#else
#define safe ((void)0)
#define debug(...) ((void)0)
#define orange(...) ((void)0)
#endif
```

1.3 Increase Stack

```
const int size = 256 << 20;</pre>
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));
```

Pragma Optimization

```
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
1.5 IO Optimization
static inline int gc() {
 constexpr int B = 1<<20;</pre>
 static char buf[B], *p, *q;
 if(p == q \&\&
  (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
  return EOF:
 return *p++;
template < typename T >
static inline bool gn( T &x ) {
 int c = gc(); T sgn = 1; x = 0;
while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
if(c == '-') sgn = -1, c = gc();
 if(c == EOF) return false;
 while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
 return x *= sgn, true;
```

#pragma GCC optimize("Ofast,no-stack-protector")

#pragma GCC optimize("no-math-errno,unroll-loops")

Data Structure 2

Dark Magic 2.1

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using namespace __gnu_pbds;
// heap tags: paring/binary/binomial/rc_binomial/thin
template<typename T>
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>, \
                  pairing_heap_tag>;
// pbds_heap::point_iterator
// x = pq.push(10); pq.modify(x, 87); a.join(b);
// tree tags: rb_tree_tag/ov_tree_tag/splay_tree_tag
template<typename T>
using ordered_set = tree<T, null_type, less<T>
   rb_tree_tag, tree_order_statistics_node_update>;
// find_by_order, order_of_key
// hash tables: cc_hash_table/gp_hash_table
```

2.2 Link-Cut Tree

p->ch[dir]=c;

```
struct Node{
Node *par, *ch[2];
int xor_sum, v;
bool is_rev;
Node(int _v){
 v=xor_sum=_v;is_rev=false;
 par=ch[0]=ch[1]=nullptr;
inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
inline void down(){
 if(is_rev){
  if(ch[0]!=nullptr) ch[0]->set_rev();
   if(ch[1]!=nullptr) ch[1]->set_rev();
   is_rev=false;
 }
inline void up(){
 xor_sum=v;
  if(ch[0]!=nullptr){
  xor_sum^=ch[0]->xor_sum;
  ch[0]->par=this;
 if(ch[1]!=nullptr){
  xor_sum^=ch[1]->xor_sum;
  ch[1]->par=this;
inline bool is_root(){
 return par==nullptr ||\
   (par->ch[0]!=this && par->ch[1]!=this);
bool is_rch(){return !is_root() && par->ch[1]==this;}
} *node[maxn], *stk[maxn];
int top;
void to_child(Node* p,Node* c,bool dir){
```

```
p->up();
inline void rotate(Node* node){
 Node* par=node->par;
 Node* par_par=par->par;
 bool dir=node->is_rch()
 bool par_dir=par->is_rch()
 to_child(par, node->ch[!dir], dir);
 to_child(node,par,!dir);
 if(par_par!=nullptr && par_par->ch[par_dir]==par)
  to_child(par_par,node,par_dir);
 else node->par=par_par;
inline void splay(Node* node){
 Node* tmp=node;
 stk[top++]=node;
 while(!tmp->is_root()){
  tmp=tmp->par;
  stk[top++]=tmp;
 while(top) stk[--top]->down();
 for(Node *fa=node->par;
  !node->is_root();
  rotate(node), fa=node->par)
  if(!fa->is_root())
   rotate(fa->is_rch()==node->is_rch()?fa:node);
inline void access(Node* node){
 Node* last=nullptr;
 while(node!=nullptr){
  splay(node);
  to_child(node, last, true);
  last=node;
  node=node->par;
inline void change_root(Node* node){
 access(node);splay(node);node->set_rev();
inline void link(Node* x, Node* y){
 change_root(x);splay(x);x->par=y;
inline void split(Node* x,Node* y){
 change_root(x);access(y);splay(x);
 to_child(x,nullptr,true);y->par=nullptr;
inline void change_val(Node* node,int v){
access(node);splay(node);node->v=v;node->up();
inline int query(Node* x, Node* y){
 change_root(x);access(y);splay(y);
 return y->xor_sum;
inline Node* find_root(Node* node){
 access(node);splay(node);
 Node* last=nullptr:
 while(node!=nullptr){
  node->down();last=node;node=node->ch[0];
 return last;
set<pii> dic;
inline void add_edge(int u,int v){
 if(u>v) swap(u,v)
 if(find_root(node[u])==find_root(node[v])) return;
 dic.insert(pii(u,v))
link(node[u],node[v]);
inline void del_edge(int u,int v){
 if(u>v) swap(u,v);
 if(dic.find(pii(u,v))==dic.end()) return;
 dic.erase(pii(u,v))
 split(node[u],node[v]);
2.3 LiChao Segment Tree
 int m, k, id;
 Line() : id( -1 ) {}
```

```
struct Line{
 Line('int a, int'b,'int c')
: m(a), k(b), id(c) {}
 int at( int x ) { return m * x + k; }
```

```
#undef sz
class LiChao {
 private:
                                                             2.5 Linear Basis
  int n; vector< Line > nodes;
  inline int lc( int x ) { return 2 * x + 1; }
                                                             template <int BITS>
  inline int rc( int x ) { return 2 * x + 2; }
                                                             struct LinearBasis {
  void insert( int 1, int r, int id, Line ln ) {
                                                              array<uint64_t, BITS> basis;
   int m = (1 + r) >> 1;
                                                              Basis() { basis.fill(0); }
                                                              void add(uint64_t x)
   if ( nodes[ id ].id == -1 ) {
    nodes[ id ] = ln;
                                                               for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
                                                                if (basis[i] == 0) {
    return:
                                                                 basis[i] = x;
   bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
                                                                 return;
   if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
    atLeft ^= 1; swap( nodes[ id ], ln );
                                                                x ^= basis[i];
                                                               }
   if ( r - 1 == 1 ) return;
   if ( atLeft ) insert( l, m, lc( id ), ln );
                                                              bool ok(uint64_t x) {
   else insert( m, r, rc( id ), ln );
                                                               for (int i = 0; i < BITS; ++i)</pre>
                                                                if ((x >> i) & 1) x ^= basis[i];
  int query( int 1, int r, int id, int x ) {
                                                               return x == 0;
   int ret = 0;
   if ( nodes[ id ].id != -1 )
                                                             };
    ret = nodes[ id ].at( x );
                                                             2.6
                                                                   Binary Search On Segment Tree
   int m = (1 + r) >> 1;
   if ( r - l == 1 ) return ret;
                                                             // find_first = x -> minimal x s.t. check( [a, x) )
   else if (x < m )
                                                             // find_last = x \rightarrow maximal x s.t. check([x, b))
    return max( ret, query( 1, m, lc( id ), x ) );
                                                             template <typename C>
   else
                                                             int find_first(int 1, const C &check) {
    return max( ret, query( m, r, rc( id ), x ) );
                                                              if (1 >= n) return n;
                                                              1 += sz;
 public:
                                                              for (int i = height; i > 0; i--)
  void build( int n_ ) {
                                                               propagate(l >> i);
  n = n_; nodes.clear();
                                                              Monoid sum = identity;
   nodes.resize( n << 2, Line() );</pre>
                                                               while ((1 & 1) == 0) 1 >>= 1;
  void insert( Line ln ) { insert( 0, n, 0, ln ); }
                                                               if (check(f(sum, data[1]))) {
  int query( int x ) { return query( 0, n, 0, x ); }
                                                                while (1 < sz) {</pre>
                                                                 propagate(1);
                                                                 1 <<= 1;
2.4 Treap
                                                                 auto nxt = f(sum, data[1]);
namespace Treap{
                                                                 if (not check(nxt)) {
 #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                                  sum = nxt;
 struct node{
                                                                  1++;
  int size;
                                                                 }
  uint32_t pri;
                                                                }
  node *lc, *rc, *pa;
                                                                return 1 + 1 - sz;
  node():size(0),pri(rand()),lc(0),rc(0),pa(0){}
  void pull() {
                                                               sum = f(sum, data[1++]);
  size = 1; pa = nullptr;
                                                              } while ((1 & -1) != 1);
   if ( lc ) { size += lc->size; lc->pa = this; }
if ( rc ) { size += rc->size; rc->pa = this; }
                                                              return n;
  }
                                                             template <typename C>
                                                             int find_last(int r, const C &check) {
node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
                                                              if (r <= 0) return -1;
                                                              r += sz;
  if ( L->pri > R->pri ) {
                                                              for (int i = height; i > 0; i--)
  L->rc = merge( L->rc, R ); L->pull();
                                                               propagate((r - 1) >> i);
   return L;
                                                              Monoid sum = identity;
  } else {
                                                              do {
   R->lc = merge( L, R->lc ); R->pull();
   return R;
                                                               while (r > 1 \text{ and } (r \& 1)) r >>= 1;
  }
                                                               if (check(f(data[r], sum))) {
                                                                while (r < sz) {</pre>
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                                 propagate(r);
  if ( not rt ) L = R = nullptr;
                                                                 r = (r << 1) + 1;
  else if( sz( rt->lc ) + 1 <= k ) {
                                                                 auto nxt = f(data[r], sum);
                                                                 if (not check(nxt)) {
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                  sum = nxt;
   L->pull();
                                                                  r--:
  } else {
                                                                 }
   R = rt;
   split_by_size( rt->lc, k, L, R->lc );
                                                                return r - sz;
   R->pull();
                                                               sum = f(data[r], sum);
  }
                                                              } while ((r & -r) != r);
 int getRank(node *o) {
                                                              return -1;
  int r = sz(o->lc);
  for (;o->pa != nullptr; o = o->pa)
   if (o->pa->rc != o) r += sz(o->pa->lc);
                                                                  Graph
  return r;
```

3.1 BCC Edge

```
class BCC_Bridge {
                                                                  for (int i = 0; i < n; ++i)
                                                                   if (not dfn[i]) dfs(i, i);
private:
 int n, ecnt;
  vector<vector<pair<int,int>>> G;
                                                                 int get_id(int x) { return bcc[x]; }
 vector<int> dfn, low;
                                                                 int count() { return ecnt; }
  vector<bool> bridge;
                                                                 bool is_ap(int x) { return ap[x]; }
 void dfs(int u, int f) {
   dfn[u] = low[u] = dfn[f] + 1;
                                                              } bcc_ap;
                                                               3.3 2-SAT (SCC)
   for (auto [v, t]: G[u]) {
  if (v == f) continue;
                                                              class TwoSat{
    if (dfn[v]) {
                                                                private:
     low[u] = min(low[u], dfn[v]);
                                                                 int n:
     continue;
                                                                 vector<vector<int>> rG,G,sccs;
                                                                 vector<int> ord,idx;
    dfs(v, u);
                                                                 vector<bool> vis,result;
   low[u] = min(low[u], low[v]);
if (low[v] > dfn[u]) bridge[t] = true;
                                                                 void dfs(int u){
                                                                  vis[u]=true
                                                                  for(int v:G[u])
  }
                                                                   if(!vis[v]) dfs(v);
public:
                                                                  ord.push_back(u);
  void init(int n_) {
   G.clear(); G.resize(n = n_);
                                                                 void rdfs(int u){
   low.assign(n, ecnt = 0);
                                                                  vis[u]=false;idx[u]=sccs.size()-1;
                                                                  sccs.back().push_back(u);
   dfn.assign(n, 0);
                                                                  for(int v:rG[u])
  void add_edge(int u, int v) {
                                                                   if(vis[v])rdfs(v);
  G[u].emplace_back(v, ecnt);
   G[v].emplace_back(u, ecnt++);
                                                                public:
                                                                 void init(int n_){
  void solve() {
                                                                  n=n_;G.clear();G.resize(n);
  bridge.assign(ecnt, false);
                                                                  rG.clear();rG.resize(n)
   for (int i = 0; i < n; ++i)
                                                                  sccs.clear();ord.clear();
    if (not dfn[i]) dfs(i, i);
                                                                  idx.resize(n);result.resize(n);
  bool is_bridge(int x) { return bridge[x]; }
                                                                 void add_edge(int u,int v){
} bcc_bridge;
                                                                  G[u].push_back(v);rG[v].push_back(u);
3.2 BCC Vertex
                                                                 void orr(int x,int y){
class BCC_AP {
                                                                  if ((x^y)==1)return
                                                                  add_edge(x^1,y); add_edge(y^1,x);
private:
 int n, ecnt;
 vector<vector<pair<int,int>>> G;
                                                                 bool solve(){
  vector<int> bcc, dfn, low, st;
                                                                  vis.clear();vis.resize(n);
  vector<bool> ap, ins;
                                                                  for(int i=0;i<n;++i)</pre>
 void dfs(int u, int f) {
  dfn[u] = low[u] = dfn[f] + 1;
                                                                   if(not vis[i])dfs(i);
                                                                  reverse(ord.begin(),ord.end());
   int ch = 0;
                                                                  for (int u:ord){
   for (auto [v, t]: G[u]) if (v != f) {
                                                                   if(!vis[u])continue:
    if (not ins[t]) {
                                                                   sccs.push_back(vector<int>());
     st.push_back(t);
                                                                   rdfs(u);
     ins[t] = true;
                                                                  for(int i=0;i<n;i+=2)</pre>
    if (dfn[v]) {
                                                                   if(idx[i]==idx[i+1])
     low[u] = min(low[u], dfn[v]);
                                                                    return false;
                                                                  vector<bool> c(sccs.size());
    } ++ch; dfs(v, u);
                                                                  for(size_t i=0;i<sccs.size();++i){</pre>
    low[u] = min(low[u], low[v]);
                                                                   for(size_t j=0;j<sccs[i].size();++j){</pre>
                                                                    result[sccs[i][j]]=c[i]
    if (low[v] >= dfn[u]) {
     ap[u] = true;
                                                                    c[idx[sccs[i][j]^1]]=!c[i];
     while (true) {
      int eid = st.back(); st.pop_back();
      bcc[eid] = ecnt;
                                                                  return true;
      if (eid == t) break;
     }
                                                                 bool get(int x){return result[x];}
                                                                 inline int get_id(int x){return idx[x];}
     ecnt++;
                                                                 inline int count(){return sccs.size();}
    }
                                                              } sat2;
   if (ch == 1 and u == f) ap[u] = false;
                                                               3.4 Lowbit Decomposition
public:
                                                              class LowbitDecomp{
  void init(int n_) {
                                                               private:
  G.clear(); G.resize(n = n_);
                                                                int time_, chain_, LOG_N;
   ecnt = 0; ap.assign(n, false);
                                                                vector< vector< int > > G, fa;
                                                                vector< int > tl, tr, chain, chain_st;
// chain_ : number of chain
   low.assign(n, 0); dfn.assign(n, 0);
                                                                // tl, tr[ u ] : subtree interval in the seq. of u
  void add_edge(int u, int v) {
                                                                // chain_st[ u ] : head of the chain contains u // chian[ u ] : chain id of the chain u is on
   G[u].emplace_back(v, ecnt);
   G[v].emplace_back(u, ecnt++);
                                                                void predfs( int u, int f ) {
  void solve() {
                                                                 chain[u] = 0;
   ins.assign(ecnt, false);
                                                                 for ( int v : G[ u ] ) {
                                                                  if ( v == f ) continue;
   bcc.resize(ecnt); ecnt = 0;
```

```
predfs( v, u );
                                                              class MaxClique{
   if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
                                                              private:
    chain[ u ] = chain[ v ];
                                                               using bits = bitset< MAXN >;
                                                               bits popped, G[ MAXN ], ans;
size_t deg[ MAXN ], deo[ MAXN ], n;
  if ( not chain[ u ] )
   chain[ u ] = chain_ ++;
                                                               void sort_by_degree() {
                                                                popped.reset();
                                                                for ( size_t i = 0 ; i < n ; ++ i )</pre>
 void dfschain( int u, int f ) {
  fa[ u ][ 0 ] = f;
for ( int i = 1 ; i < LOG_N ; ++ i )
                                                                  deg[ i ] = G[ i ].count();
                                                                for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;</pre>
   fa[u][i] = fa[fa[u][i-1]][i-1];
                                                                  for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )
    mi = deg[ id = j ];</pre>
  tl[ u ] = time_++;
  if ( not chain_st[ chain[ u ] ] )
   chain_st[ chain[ u ] ] = u;
                                                                  popped[ deo[ i ] = id ] = 1;
  for ( int v : G[ u ] )
   if ( v != f and chain[ v ] == chain[ u ] )
                                                                  for( size_t u = G[ i ]._Find_first() ;
  dfschain( v, u );
for ( int v : G[ u ] )
                                                                   u < n ; u = G[ i ]._Find_next( u ) )</pre>
                                                                     -- deg[ u ];
   if ( v != f and chain[ v ] != chain[ u ] )
                                                                }
    dfschain( v, u );
                                                               void BK( bits R, bits P, bits X ) {
  tr[ u ] = time_;
                                                                if (R.count()+P.count() <= ans.count()) return;</pre>
                                                                if ( not P.count() and not X.count() ) {
 bool anc( int u, int v )
 return tl[ u ] <= tl[ v ] and tr[ v ] <= tr[ u ];</pre>
                                                                 if ( R.count() > ans.count() ) ans = R;
                                                                 return:
public:
                                                                }
                                                                /* greedily chosse max degree as pivot
 int lca( int u, int v ) {
  if ( anc( u, v ) ) return u;
                                                                bits cur = P | X; size_t pivot = 0, sz = 0;
  for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
                                                                for ( size_t u = cur._Find_first() ;
   if ( not anc( fa[ u ][ i ], v ) )
                                                                 u < n ; u = cur._Find_next( u )
    u = fa[ u ][ i ];
                                                                  if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
                                                                cur = P & ( ~G[ pivot ] );
  return fa[ u ][ 0 ];
                                                                */ // or simply choose first
                                                                bits cur = P & (~G[ ( P | X )._Find_first() ]);
 void init( int n ) {
  fa.assign( ++n, vector< int >( LOG_N ) );
                                                                for ( size_t u = cur._Find_first()
  for (LOG_N = 0 ; (1 << LOG_N ) < n ; ++ LOG_N );
                                                                 u < n ; u = cur._Find_next( u ) ) {
                                                                 if ( R[ u ] ) continue;
  G.clear(); G.resize( n );
  tl.assign( n, 0 ); tr.assign( n, 0 )
                                                                 R[u] = 1;
  chain.assig( n, 0 ); chain_st.assign( n, 0 );
                                                                 BK( R, P & G[ u ], X & G[ u ] );
                                                                 R[u] = P[u] = 0, X[u] = 1;
 void add_edge( int u , int v ) {
  // 1-base
  G[ u ].push_back( v );
                                                              public:
  G[ v ].push_back( u );
                                                               void init( size_t n_ ) {
                                                                n = n_{-};
 }
 void decompose(){
                                                                for ( size_t i = 0 ; i < n ; ++ i )
                                                                 G[ i ].reset();
 chain_ = 1;
 predfs( 1, 1 );
                                                                ans.reset();
  time_{-} = 0;
 dfschain( 1, 1 );
                                                               void add_edges( int u, bits S ) { G[ u ] = S; }
                                                               void add_edge( int u, int v ) {
 PII get_subtree(int u) { return {tl[ u ],tr[ u ] }; }
                                                                G[u][v] = G[v][u] = 1;
 vector< PII > get_path( int u , int v ){
  vector< PII > res;
                                                               int solve() {
  int g = lca( u, v );
                                                                sort_by_degree(); // or simply iota( deo... )
  while ( chain[ u ] != chain[ g ] ) {
                                                                for ( size_t i = 0 ; i < n ; ++ i )</pre>
   int s = chain_st[ chain[ u ] ];
                                                                 deg[ i ] = G[ i ].count();
   res.emplace_back( tl[ s ], tl[ u ] + 1 );
                                                                bits pob, nob = 0; pob.set();
   u = fa[ s ][ 0 ];
                                                                for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
                                                                for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
  res.emplace_back( tl[ g ], tl[ u ] + 1 );
while ( chain[ v ] != chain[ g ] ) {
                                                                 size_t v = deo[ i ];
                                                                 bits tmp; tmp[ v ] = 1;
  int s = chain_st[ chain[ v ] ];
                                                                 BK( tmp, pob & G[ v ], nob & G[ v ] );
   res.emplace_back( tl[ s ], tl[ v ] + 1 );
                                                                 pob[v] = 0, nob[v] = 1;
   v = fa[ s ][ 0 ];
                                                                return static_cast< int >( ans.count() );
  res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
                                                               }
  return res;
                                                              };
  /* res : list of intervals from u to v
                                                                   MaxCliqueDyn
   \star ( note only nodes work, not edge )
                                                              constexpr int kN = 150;
   * vector< PII >& path = tree.get_path( u , v )
                                                              struct MaxClique { // Maximum Clique
                                                               bitset<kN> a[kN], cs[kN];
   * for( auto [ 1, r ] : path ) {
   * 0-base [ 1, r )
                                                               int ans, sol[kN], q, cur[kN], d[kN], n;
   * }
                                                               void init(int _n) {
   */
                                                               n = n, ans q = 0;
                                                                for (int i = 0; i < n; i++) a[i].reset();</pre>
} tree;
                                                               void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
3.5 MaxClique
                                                               void csort(vector<int> &r, vector<int> &c) {
// contain a self loop u to u, than u won't in clique
                                                                int mx = 1, km = max(ans - q + 1, 1), t = 0,
                                                                  m = int(r.size());
template < size_t MAXN >
```

```
cs[1].reset(); cs[2].reset();
                                                                 addEdge(stk[i], stk[i + 1]);
  for (int i = 0; i < m; i++) {
  int p = r[i], k = 1;</pre>
                                                               3.8 Centroid Decomposition
   while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
                                                               struct Centroid {
   cs[k][p] = 1;
                                                                vector<vector<int64_t>> Dist;
   if (k < km) r[t++] = p;
                                                                vector<int> Parent, Depth;
                                                                vector<int64_t> Sub, Sub2;
  c.resize(m);
                                                                vector<int> Sz, Sz2;
  if (t) c[t - 1] = 0;
                                                                Centroid(vector<vector<pair<int, int>>> g) {
  for (int k = km; k <= mx; k++) {</pre>
                                                                 int N = g.size()
  for (int p = int(cs[k]._Find_first());
                                                                 vector<bool> Vis(N);
      p < kN; p = int(cs[k]._Find_next(p))) {</pre>
                                                                 vector<int> sz(N), mx(N);
    r[t] = p; c[t++] = k;
                                                                 vector<int> Path;
                                                                 Dist.resize(N)
                                                                 Parent.resize(N);
                                                                 Depth.resize(N)
 void dfs(vector<int> &r, vector<int> &c, int 1,
                                                                 auto DfsSz = [\&](auto dfs, int x) -> void {
  bitset<kN> mask) {
                                                                  Vis[x] = true; sz[x] = 1; mx[x] = 0;
  while (!r.empty()) {
                                                                  for (auto [u, w] : g[x]) {
                                                                   if (Vis[u]) continue;
   int p = r.back(); r.pop_back();
   mask[p] = 0;
                                                                   dfs(dfs, u)
   if (q + c.back() <= ans) return;</pre>
                                                                   sz[x] += sz[u];
   cur[q++] = p;
                                                                   mx[x] = max(mx[x], sz[u]);
   vector<int> nr, nc;
   bitset<kN> nmask = mask & a[p];
                                                                  Path.push_back(x);
   for (int i : r)
                                                                 };
    if (a[p][i]) nr.push_back(i);
                                                                 auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
   if (!nr.empty()) {
                                                                  -> void {
                                                                  Dist[x].push_back(D);Vis[x] = true;
    if (1 < 4) {
     for (int i : nr)
                                                                  for (auto [u, w] : g[x]) {
                                                                   if (Vis[u]) continue;
      d[i] = int((a[i] & nmask).count());
     sort(nr.begin(), nr.end(),
                                                                   dfs(dfs, u, D + w);
      [&](int x, int y)
       return d[x] > d[y];
      });
                                                                 auto Dfs = [&]
                                                                  (auto dfs, int x, int D = 0, int p = -1)->void {
   csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
                                                                  Path.clear(); DfsSz(DfsSz, x);
                                                                  int M = Path.size();
    ans = q; copy(cur, cur + q, sol);
                                                                  int C = -1:
                                                                  for (int u : Path) {
   c.pop_back(); q--;
                                                                   if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
  }
                                                                   Vis[u] = false;
 int solve(bitset<kN> mask) { // vertex mask
                                                                  DfsDist(DfsDist, C);
                                                                  for (int u : Path) Vis[u] = false;
  vector<int> r, c;
  for (int i = 0; i < n; i++)
                                                                  Parent[C] = p; Vis[C] = true;
  if (mask[i]) r.push_back(i);
for (int i = 0; i < n; i++)</pre>
                                                                  Depth[C] = D;
                                                                  for (auto [u, w] : g[C]) {
  d[i] = int((a[i] & mask).count());
                                                                   if (Vis[u]) continue;
  sort(r.begin(), r.end(),
                                                                   dfs(dfs, u, D + 1, C);
   [&](int i, int j) { return d[i] > d[j]; });
  csort(r, c);
  dfs(r, c, 1, mask);
                                                                 Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
  return ans; // sol[0 ~ ans-1]
                                                                 Sz.resize(N); Sz2.resize(N);
} graph;
                                                                void Mark(int v) {
                                                                 int x = v, z = -1
3.7 Virtural Tree
                                                                 for (int i = Depth[v]; i >= 0; --i) {
                                                                  Sub[x] += Dist[v][i]; Sz[x]++;
inline bool cmp(const int &i, const int &j) {
                                                                  if (z != -1) {
return dfn[i] < dfn[j];</pre>
                                                                   Sub2[z] += Dist[v][i];
void build(int vectrices[], int k) {
                                                                   Sz2[z]++;
 static int stk[MAX_N];
 sort(vectrices, vectrices + k, cmp);
                                                                  z = x; x = Parent[x];
 stk[sz++] = 0;
                                                                 }
 for (int i = 0; i < k; ++i) {
  int u = vectrices[i], lca = LCA(u, stk[sz - 1]);
  if (lca == stk[sz - 1]) stk[sz++] = u;</pre>
                                                                int64_t Query(int v) {
                                                                 int64_t res = 0;
                                                                 int x = v, z = -1
                                                                 for (int i = Depth[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
   while (sz \ge 2 \&\& dep[stk[sz - 2]] \ge dep[lca]) {
    addEdge(stk[sz - 2], stk[sz - 1]);
                                                                  if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
                                                                  z = x; x = Parent[x];
   if (stk[sz - 1] != lca) {
   addEdge(lca, stk[--sz]);
                                                                 return res;
    stk[sz++] = lca, vectrices[cnt++] = lca;
                                                               };
   stk[sz++] = u;
                                                               3.9 Tree Hashing
  }
                                                              |uint64_t hsah(int u, int f) {
 for (int i = 0; i < sz - 1; ++i)
                                                              uint64_t r = 127;
```

```
for (int v : G[ u ]) if (v != f) {
  uint64_t hh = hsah(v, u);
                                                                        stk[ stk_ ++ ] = u;
  r=(r+(hh*hh)%1010101333)%1011820613;
                                                                       bool inPath[ N ];
                                                                       void Diff( int u ) {
return r:
                                                                        if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
}
                                                                        else { /*add this edge*/ }
3.10 Minimum Mean Cycle
/* minimum mean cycle O(VE) */
                                                                       void traverse( int& origin_u, int u ) {
                                                                        for ( int g = lca( origin_u, u )
struct MMC{
                                                                         origin_u != g ; origin_u = parent_of[ origin_u ] )
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
                                                                          Diff( origin_u );
#define V 1021
                                                                        for (int v = u; v != origin_u; v = parent_of[v])
#define inf 1e9
                                                                         Diff( v );
 struct Edge { int v,u; double c; };
                                                                        origin_u = u;
 int n, m, prv[V][V], prve[V][V], vst[V];
                                                                       void solve() {
 Edge e[E];
                                                                        dfs( 1, 1 );
 vector<int> edgeID, cycle, rho;
 double d[V][V];
                                                                        while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
                                                                        sort( que, que + q, [](const Que& x, const Que& y) {
  return tie( block_id[ x.u ], dfn[ x.v ] )
 void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
 void add_edge( int vi , int ui , double ci )
                                                                              < tie( block_id[ y.u ], dfn[ y.v ] );
 { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
                                                                        } );
                                                                        int U = 1, V = 1;
  for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);</pre>
                                                                        for ( int i = 0 ; i < q ; ++ i ) {
  pass( U, que[ i ].u );
  pass( V, que[ i ].v );</pre>
   for(int j=0; j<m; j++) {</pre>
                                                                         // we could get our answer of que[ i ].id
    int v = e[j].v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
                                                                      }
                                                                       /*
     d[i+1][u] = d[i][v]+e[j].c;
      prv[i+1][u] = v;
                                                                      Method 2:
      prve[i+1][u] = j;
                                                                      dfs u:
                                                                       push u
                                                                        iterate subtree
                                                                       Let P = LCA(u, v), and St(u) \le St(v)
                                                                      if (P == u) query[St(u), St(v)]
 double solve(){
  // returns inf if no cycle, mmc otherwise
                                                                       else query[Ed(u), St(v)], query[St(P), St(P)]
  double mmc=inf;
  int st = -1;
                                                                       3.12 Minimum Steiner Tree
  bellman_ford();
  for(int i=0; i<n; i++) {</pre>
                                                                      // Minimum Steiner Tree
   double avg=-inf;
                                                                       // 0(V 3^T + V^2 2^T)
                                                                      struct SteinerTree{
   for(int k=0; k<n; k++) {</pre>
    if(d[n][i]<inf-eps)</pre>
                                                                       #define V 33
      avg=max(avg,(d[n][i]-d[k][i])/(n-k));
                                                                       #define T 8
                                                                       #define INF 1023456789
     else avg=max(avg,inf);
                                                                        int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
                                                                        void init( int _n ){
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
                                                                         n = _n;
                                                                         for( int i = 0 ; i < n ; i ++ ){
  for( int j = 0 ; j < n ; j ++ )
  dst[ i ][ j ] = INF;</pre>
  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
  for (int i=n; !vst[st]; st=prv[i--][st]) {
   vst[st]++
   edgeID.PB(prve[i][st]);
                                                                          dst[ i ][ i ] = 0;
   rho.PB(st);
  while (vst[st] != 2) {
  int v = rho.back(); rho.pop_back();
                                                                        void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
   cycle.PB(v);
                                                                         dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
   vst[v]++;
                                                                        void shortest_path(){
  reverse(ALL(edgeID));
                                                                         for( int k = 0 ; k < n ; k ++ )</pre>
                                                                          for( int i = 0 ; i < n ; i ++ )
for( int j = 0 ; j < n ; j ++ )
dst[ i ][ j ] = min( dst[ i ][ j ],
  edgeID.resize(SZ(cycle));
  return mmc;
} mmc;
                                                                                 dst[ i ][ k ] + dst[ k ][ j ] );
3.11 Mo's Algorithm on Tree
                                                                        int solve( const vector<int>& ter ){
                                                                         int t = (int)ter.size();
int q; vector< int > G[N];
                                                                         for( int i = 0 ; i < (1 << t) ; i ++ )
for( int j = 0 ; j < n ; j ++ )
struct Que{
 int u, v, id;
                                                                            dp[ i ][ j ] = INF;
} que[ N ];
int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
void dfs( int u, int f ) {
                                                                         for( int i = 0 ; i < n ; i ++ )</pre>
                                                                          dp[0][i] = 0;
 dfn[ u ] = dfn_++; int saved_rbp = stk_;
                                                                         for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
 for ( int v : G[ u ] ) {
  if ( v == f ) continue;
                                                                          if( msk == ( msk & (-msk) ) ){
                                                                           int who = __lg( msk );
  dfs( v, u );
                                                                            for( int i = 0 ; i < n ; i ++ )</pre>
  if ( stk_ - saved_rbp < SQRT_N ) continue;
for ( ++ block_ ; stk_ != saved_rbp ; )
  block_id[ stk[ -- stk_ ] ] = block_;</pre>
                                                                             dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
                                                                            continue;
```

```
National Taiwan University - ckiseki
   for( int i = 0 ; i < n ; i ++ )</pre>
    for( int submsk = ( msk - 1 ) & msk ; submsk ;
          submsk = ( submsk - 1 ) & msk )
       dp[ msk ^ submsk ][ i ] );
   for( int i = 0 ; i < n ; i ++ ){</pre>
    tdst[ i ] = INF;
    for( int j = 0 ; j < n ; j ++ )
tdst[ i ] = min( tdst[ i ],
    dp[ msk ][ j ] + dst[ j ][ i ] );</pre>
   for( int i = 0 ; i < n ; i ++ )
dp[ msk ][ i ] = tdst[ i ];</pre>
  int ans = INF;
  for( int i = 0 ; i < n ; i ++ )
ans = min( ans , dp[ ( 1 << t ) - 1 ][ i ] );
  return ans;
} solver;
      Directed Minimum Spanning Tree
struct DirectedMST { // find maximum
 struct Edge {
  int u, v;
  int w;
  Edge(int u, int v, int w) : u(u), v(v), w(w) {}
 vector<Edge> Edges;
 void clear() { Edges.clear(); }
```

```
void addEdge(int a, int b, int w) { Edges.emplace_back
 (a, b, w); }
int solve(int root, int n) {
  vector<Edge> E = Edges;
  int ans = 0;
  while (true) {
   // find best in edge
   vector<int> in(n, -inf), prv(n, -1);
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
     in[e.v] = e.w;
     prv[e.v] = e.u;
   in[root] = 0;
   prv[root] = -1;
   for (int i = 0; i < n; i++)
    if (in[i] == -inf)
     return -inf;
   // find cycle
   int tot = 0;
   vector<int> id(n, -1), vis(n, -1);
for (int i = 0; i < n; i++) {</pre>
    ans += in[i];
    for (int x = i; x != -1 && id[x] == -1; x = prv[x])
     if (vis[x] == i) {
      for (int y = prv[x]; y != x; y = prv[y])
       id[y] = tot;
      id[x] = tot++;
      break;
     vis[x] = i;
    }
   if (!tot)
    return ans;
   for (int i = 0; i < n; i++)</pre>
    if (id[i] == -1)
     id[i] = tot++;
   // shrink
   for (auto &e : E) {
    if (id[e.u] != id[e.v])
     e.w -= in[e.v];
    e.u = id[e.u], e.v = id[e.v];
   n = tot:
   root = id[root];
  assert(false);
} DMST;
```

3.14 Manhattan Minimum Spanning Tree

```
typedef Point<int> P;
vector<array<int, 3>> manhattanMST(vector<P> ps) {
 vi id(sz(ps));
 iota(all(id), 0);
 vector<array<int, 3>> edges;
 rep(k, 0, 4) {
  sort(all(id), [&](int i, int j) {
   return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
  });
  map<int, int> sweep;
  for (int i : id) {
   for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++)) {
    int j = it->second;
    P d = ps[i] - ps[j];
    if (d.y > d.x) break;
    edges.push_back(\{d.y + d.x, i, j\});
   sweep[-ps[i].y] = i;
  for (P &p : ps)
   if (k \& 1) p.x = -p.x;
   else swap(p.x, p.y);
 return edges; // [{w, i, j}, ...]
}
```

```
Dominator Tree
3.15
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
 // vertices are numbered from 0 to n-1
 fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
 fill(fa, fa + n, -1); fill(val, val + n, -1);
 fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
 fill(dom, dom + n, -1); tk = 0;
 for (int i = 0; i < n; ++i) {
  g[i].clear(); r[i].clear(); rdom[i].clear();
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
 rev[dfn[x] = tk] = x;
 fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
 for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
  r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
 if (fa[x] == x) return c ? -1 : x;
 int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
 if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
 fa[x] = p;
 return c ? p : val[x];
vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
// p[i] = -2 if i is unreachable from s
 dfs(s);
 for (int i = tk - 1; i >= 0; --i) {
  for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
  if (i) rdom[sdom[i]].push_back(i);
  for (int &u : rdom[i]) {
   int p = find(u);
   if (sdom[p] == i) dom[u] = i;
   else dom[u] = p;
  if (i) merge(i, rp[i]);
 vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)</pre>
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
 for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
 return p;
}}
```

3.16 Edge Coloring

```
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
for (int i = 0; i <= N; i++)
 for (int j = 0; j <= N; j++)
C[i][j] = G[i][j] = 0;</pre>
void solve(vector<pair<int, int>> &E, int N) {
int X[kN] = {}, a;
auto update = [&](int u)
 for (X[u] = 1; C[u][X[u]]; X[u]++);
auto color = [&](int u, int v, int c) {
 int p = G[u][v];
G[u][v] = G[v][u] = c;
 C[u][c] = v, C[v][c] = u;
 C[u][p] = C[v][p] = 0;
  if(p) X[u] = X[v] = p
 else update(u), update(v);
  return p;
 };
auto flip = [&](int u, int c1, int c2) {
 int p = C[u][c1];
 swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
 if (!C[u][c1]) X[u] = c1;
 if (!C[u][c2]) X[u] = c2;
  return p;
for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {
 auto [u, v] = E[t];
  int v0 = v, c = X[u], c0 = c, d;
  vector<pair<int,
                    int>> L; int vst[kN] = {};
 while (!G[u][v0]) {
   L.emplace_back(v, d = X[v]);
   if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
     c = color(u, L[a].first, c);
   else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
     color(u, L[a].first, L[a].second);
   else if (vst[d]) break
   else vst[d] = 1, v = C[u][d];
 if (!G[u][v0]) {
   for (; v; v = flip(v, c, d), swap(c, d));
   if (C[u][c0]) { a = int(L.size()) - 1;
    while (--a >= 0 \&\& L[a].second != c)
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
   } else t--;
```

4 Matching & Flow

4.1 Kuhn Munkres

```
class KM {
private:
static constexpr 1ld INF = 1LL << 60;</pre>
vector<lld> hl,hr,slk;
vector<int> fl,fr,pre,qu;
vector<vector<lld>> w;
vector<bool> v1,vr;
int n, ql, qr;
bool check(int x) {
 if (vl[x] = true, fl[x] != -1)
  return vr[qu[qr++] = fl[x]] = true;
 while (x != -1) swap(x, fr[fl[x] = pre[x]]);
 return false;
 void bfs(int s) {
 fill(slk.begin(), slk.end(), INF);
 fill(vl.begin(), vl.end(), false)
  fill(vr.begin(), vr.end(), false);
  q1 = qr = 0;
  vr[qu[qr++] = s] = true;
  while (true) {
   11d d;
   while (ql < qr) {</pre>
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
```

```
if(!vl[x]&&slk[x]>=(d=hl[x]+hr[y]-w[x][y])){
       if (pre[x] = y, d) slk[x] = d;
       else if (!check(x)) return;
    }
   d = INF;
   for (int x = 0; x < n; ++x)
    if (!v1[x] && d > s1k[x]) d = s1k[x];
   for (int x = 0; x < n; ++x) {
  if (v1[x]) h1[x] += d;
    else slk[x] -= d;
    if (vr[x]) hr[x] -= d;
   for (int x = 0; x < n; ++x)
    if (!v1[x] && !s1k[x] && !check(x)) return;
public:
 void init( int n_ ) {
  qu.resize(n = n_);
  fl.assign(n, -1); fr.assign(n, -1);
  hr.assign(n, 0); hl.resize(n);
w.assign(n, vector<lld>(n));
  slk.resize(n); pre.resize(n);
  vl.resize(n); vr.resize(n);
 void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 11d solve() {
  for (int i = 0; i < n; ++i)</pre>
   hl[i] = *max_element(w[i].begin(), w[i].end());
  for (int i = 0; i < n; ++i) bfs(i);
  11d res = 0;
  for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
  return res;
 }
} km;
4.2 Bipartite Matchina
struct BipartiteMatching {
 vector<int> X[N];
 int fX[N], fY[N], n;
 bitset<N> vis;
 bool dfs(int x)
  for (auto i:X[x]) {
  if (vis[i]) continue;
   vis[i] = true;
   if (fY[i]==-1 || dfs(fY[i])){
    fY[fX[x] = i] = x;
    return true:
   }
  return false;
 void init(int n_, int m) {
  vis.reset();
  fill(X, X + (n = n_), vector<int>());
memset(fX, -1, sizeof(int) * n);
memset(fY, -1, sizeof(int) * m);
 void add_edge(int x, int y){
  X[x].push_back(y); }
 int solve() { // return how many pair matched
  int cnt = 0;
  for(int i=0;i<n;i++) {</pre>
   vis.reset()
   cnt += dfs(i);
  return cnt;
4.3 General Graph Matching
namespace matching {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
 for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
```

for (int i = 0; i < n; ++i) g[i].clear();</pre>

void AddEdge(int u, int v) {

dis[m] = dis[u] - edge[v][m] + edge[u][v];

```
g[u].push_back(v);
                                                                    onstk[v] = 1;
g[v].push_back(u);
                                                                    stk.PB(v);
                                                                    if (SPFA(m)) return true;
int Find(int u) {
                                                                    stk.pop_back();
return u == fa[u] ? u : fa[u] = Find(fa[u]);
                                                                    onstk[v] = 0;
int LCA(int x, int y, int n) {
static int tk = 0; tk++;
                                                                 onstk[u] = 0; stk.pop_back();
x = Find(x), y = Find(y);
for (; ; swap(x, y)) {
  if (x != n) {
                                                                 return false;
  if (v[x] == tk) return x;
                                                                int solve() { // find a match
                                                                for (int i=0; i<n; i+=2){
  match[i] = i+1;</pre>
  v[x] = tk;
   x = Find(pre[match[x]]);
                                                                  match[i+1] = i;
                                                                 while (true){
void Blossom(int x, int y, int 1) {
  while (Find(x) != 1) {
                                                                  int found = 0;
                                                                  for (int i=0; i<n; i++)</pre>
 pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
                                                                   dis[i] = onstk[i] = 0;
                                                                  for (int i=0; i<n; i++){</pre>
  if (fa[x] == x) fa[x] = 1;
                                                                   stk.clear()
  if (fa[y] == y) fa[y] = 1;
                                                                   if (!onstk[i] && SPFA(i)){
                                                                    found = 1
 x = pre[y];
                                                                    while (SZ(stk)>=2){
                                                                     int u = stk.back(); stk.pop_back();
                                                                     int v = stk.back(); stk.pop_back();
bool Bfs(int r, int n) {
for (int i = 0; i \le n; ++i) fa[i] = i, s[i] = -1;
                                                                     match[u] = v;
while (!q.empty()) q.pop();
                                                                     match[v] = u;
q.push(r);
s[r] = 0;
                                                                   }
while (!q.empty()) {
  int x = q.front(); q.pop();
                                                                  if (!found) break;
  for (int u : g[x]) {
  if (s[u] == -1) {
                                                                 int ret = 0;
    pre[u] = x, s[u] = 1;
                                                                 for (int i=0; i<n; i++)</pre>
    if (match[u] == n) {
                                                                 ret += edge[i][match[i]];
     for (int a = u, b = x, last; b != n; a = last, b =
                                                                 return ret>>1;
     pre[a])
      last = match[b], match[b] = a, match[a] = b;
                                                              } graph;
     return true;
                                                               4.5 Minimum Cost Circulation
    q.push(match[u]);
                                                              struct Edge { int to, cap, rev, cost; };
    s[match[u]] = 0;
                                                              vector<Edge> g[kN];
                                                              int dist[kN], pv[kN], ed[kN];
   } else if (!s[u] && Find(u) != Find(x)) {
    int 1 = LCA(u, x, n);
                                                              bool mark[kN];
    Blossom(x, u, 1);
                                                              int NegativeCycle(int n) {
   Blossom(u, x, 1);
                                                               memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
  }
                                                                int upd = -1;
}
                                                                for (int i = 0; i <= n; ++i)
                                                                for (int j = 0; j < n; ++j) {
return false;
                                                                  int idx = 0:
int Solve(int n) {
                                                                  for (auto &e : g[j]) {
                                                                   if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
int res = 0;
for (int x = 0; x < n; ++x) {
                                                                    dist[e.to] = dist[j] + e.cost;
 if (match[x] == n) res += Bfs(x, n);
                                                                    pv[e.to] = j, ed[e.to] = idx;
}
                                                                    if (i == n) {
return res;
                                                                     upd = j;
                                                                     while(!mark[upd])mark[upd]=1,upd=pv[upd];
                                                                     return upd;
      Minimum Weight Matching (Clique version)
                                                                    }
struct Graph {
                                                                   idx++;
// 0-base (Perfect Match)
 int n, edge[MXN][MXN];
int match[MXN], dis[MXN], onstk[MXN];
                                                                 }
vector<int> stk;
void init(int _n) {
                                                                return -1;
 for (int i=0; i<n; i++) for (int j=0; j<n; j++)
                                                               int Solve(int n) {
   edge[i][j] = 0;
                                                               int rt = -1, ans = 0;
                                                                while ((rt = NegativeCycle(n)) >= 0) {
void set_edge(int u, int v, int w) {
                                                                 memset(mark, false, sizeof(mark));
                                                                vector<pair<int, int>> cyc;
while (!mark[rt]) {
 edge[u][v] = edge[v][u] = w; }
 bool SPFA(int u){
 if (onstk[u]) return true;
                                                                  cyc.emplace_back(pv[rt], ed[rt]);
                                                                  mark[rt] = true;
  stk.PB(u); onstk[u] = 1;
  for (int v=0; v<n; v++){
                                                                  rt = pv[rt];
  if (u != v && match[u] != v && !onstk[v]){
    int m = match[v]
                                                                 reverse(cyc.begin(), cyc.end());
    if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
                                                                 int cap = kInf;
```

for (auto &i : cyc) {

```
auto &e = g[i.first][i.second];
  cap = min(cap, e.cap);
}
for (auto &i : cyc) {
  auto &e = g[i.first][i.second];
  e.cap -= cap;
  g[e.to][e.rev].cap += cap;
  ans += e.cost * cap;
}
return ans;
}
```

4.6 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 To minimize, let f be the maximum flow from S to T. Connect
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching ${\cal M}$ on bipartite graph (X,Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect $x \to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y \to x$ with (cost,cap)=(-c,1)
 - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost, cap)=(0,d(v))
 - 5. For each vertex v with d(v)<0 , connect $v\to T$ with (cost,cap)=(0,-d(v))
 - 6. Flow from S to T , the answer is the cost of the flow $C+{\cal K}$
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \to v$, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G , connect $u\to v$ and $v\to u$ with capacity w
 - 5. For $v\in G$, connect it with sink $v\to t$ with capacity $K+2T-(\sum_{e\in E(v)}w(e))-2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight v(u,v)
 - 2. Connect v o v' with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with capacity $c_y. \\$
- 2. Create edge (x,y) with capacity c_{xy} .
- 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4.7 Dinic

```
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
    int to, rev;
    Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
       int u = bfs.front(); bfs.pop();
       for (auto e: G[u]) {
         if (e.cap <= 0 or lv[e.to]!=-1) continue;
         bfs.push(e.to); lv[e.to] = lv[u] + 1;
      }
    }
    return lv[ed] != -1;
  Cap DFS(int u, Cap f){
  if (u == ed) return f;
    Cap ret = 0;
    for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
       auto &e = G[u][i];
       if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
      Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
      G[e.to][e.rev].cap += nf;
      if (f == 0) return ret;
    if (ret == 0) lv[u] = -1;
    return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
    G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
    st = st_, ed = ed_; Cap ret = 0;
    while (BFS()) {
      idx.assign(n, 0);
      Cap f = DFS(st, numeric_limits<Cap>::max());
       ret += f;
       if (f == 0) break;
     return ret;
```

4.8 Minimum Cost Maximum Flow

```
class MiniCostMaxiFlow{
 using Cap = int; using Wei = int64_t;
 using PCW = pair<Cap,Wei>;
 static constexpr Cap INF_CAP = 1 << 30;</pre>
 static constexpr Wei INF_WEI = 1LL<<60;</pre>
private:
 struct Edge{
  int to, back;
  Cap cap; Wei wei;
  Edge() {}
  Edge(int a,int b, Cap c, Wei d):
   to(a),back(b),cap(c),wei(d) {}
 int ori, edd;
 vector<vector<Edge>> G;
 vector<int> fa, wh;
 vector<bool> inq;
 vector<Wei> dis;
 PCW SPFA(){
  fill(inq.begin(),inq.end(),false);
  fill(dis.begin(), dis.end(), INF_WEI);
  queue<int> qq; qq.push(ori);
  dis[ori] = 0;
  while(not qq.empty()){
```

```
int u=qq.front();qq.pop();
                                                                   if (c == -1 \mid | g[i] > g[c]) c = i;
   ing[u] = false;
                                                                 if (c == -1) break;
   for(int i=0;i<SZ(G[u]);++i){</pre>
                                                                 v[s = t, t = c] = true;
for (int i = 0; i < n; ++i) {
    Edge e=G[u][i];
    int v=e.to; Wei d=e.wei;
                                                                   if (del[i] | v[i]) continue;
    if(e.cap <= 0 | |dis[v] <= dis[u] + d)
                                                                   g[i] += w[c][i];
     continue;
    dis[v] = dis[u] + d;
    fa[v] = u, wh[v] = i;
    if (inq[v]) continue;
                                                                 return make_pair(s, t);
    qq.push(v);
    inq[v] = true;
                                                                int mincut(int n) {
                                                                int cut = 1e9;
  }
                                                                 memset(del, false, sizeof(del));
  if(dis[edd]==INF_WEI) return {-1, -1};
                                                                 for (int i = 0; i < n - 1; ++i) {
                                                                 int s, t; tie(s, t) = phase(n);
del[t] = true; cut = min(cut, g[t]);
for (int j = 0; j < n; ++j) {</pre>
 Cap mw=INF_CAP;
  for(int i=edd;i!=ori;i=fa[i])
  mw=min(mw,G[fa[i]][wh[i]].cap);
                                                                   w[s][j] += w[t][j]; w[j][s] += w[j][t];
 for (int i=edd;i!=ori;i=fa[i]){
   auto &eg=G[fa[i]][wh[i]];
   eg.cap -= mw;
                                                                 }
   G[eg.to][eg.back].cap+=mw;
                                                                 return cut;
  return {mw, dis[edd]};
                                                                4.11 Dijkstra Cost Flow
public:
                                                               // kN = #(vertices)
void init(int n){
                                                               // MCMF.{Init, AddEdge, MincostMaxflow}
  G.clear();G.resize(n);
                                                               // MincostMaxflow(source, sink, flow_limit, &cost)
  fa.resize(n);wh.resize(n);
                                                               // => flow
  inq.resize(n); dis.resize(n);
                                                               using Pii = pair<int, int>;
                                                               constexpr int kInf = 0x3f3f3f3f, kN = 500;
void add_edge(int st, int ed, Cap c, Wei w){
G[st].emplace_back(ed,SZ(G[ed]),c,w);
                                                               struct Edge {
                                                                int to, rev, cost, flow;
 G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
                                                               }:
                                                               struct MCMF { // 0-based
                                                                int n{}, m{}, s{}, t{};
vector<Edge> graph[kN];
PCW solve(int a, int b){
 ori = a, edd = b;
 Cap cc=0; Wei ww=0;
                                                                 // Larger range for relabeling
 while(true){
                                                                 int64_t dis[kN] = {}, h[kN] = {};
  PCW ret=SPFA();
                                                                 int p[kN] = {};
  if(ret.first==-1) break;
                                                                 void Init(int nn) {
   cc+=ret.first;
                                                                 n = nn;
  ww+=ret.first * ret.second;
                                                                 for (int i = 0; i < n; i++) graph[i].clear();</pre>
                                                                 void AddEdge(int u, int v, int f, int c) {
 return {cc,ww};
                                                                 graph[u].push_back({v
} mcmf;
                                                                   static_cast<int>(graph[v].size()), c, f});
                                                                  graph[v].push_back(
4.9 GomoryHu Tree
                                                                   {u, static_cast<int>(graph[u].size()) - 1,
int g[maxn];
                                                                    -c, 0});
vector<edge> GomoryHu(int n){
 vector<edge> rt;
                                                                 bool Dijkstra(int &max_flow, int64_t &cost) {
for(int i=1;i<=n;++i)g[i]=1;</pre>
                                                                 priority_queue<Pii, vector<Pii>, greater<>> pq;
for(int i=2;i<=n;++i){</pre>
                                                                  fill_n(dis, n, kInf);
  int t=g[i];
                                                                  dis[s] = 0;
 flow.reset(); // clear flows on all edge
                                                                  pq.emplace(0, s);
  rt.push_back({i,t,flow(i,t)});
                                                                  while (!pq.empty()) {
 flow.walk(i); // bfs points that connected to i (use
  edges not fully flow)
                                                                   auto u = pq.top();
                                                                   pq.pop();
 for(int j=i+1;j<=n;++j){</pre>
                                                                   int v = u.second;
  if(g[j]==t && flow.connect(j))g[j]=i; // check if i
                                                                   if (dis[v] < u.first) continue;</pre>
                                                                   for (auto &e : graph[v]) {
    can reach i
                                                                    auto new_dis =
                                                                     dis[v] + e.cost + h[v] - h[e.to];
return rt;
                                                                    if (e.flow > 0 && dis[e.to] > new_dis) {
                                                                     dis[e.to] = new_dis;
                                                                     p[e.to] = e.rev;
4.10 Global Min-Cut
                                                                     pq.emplace(dis[e.to], e.to);
const int maxn = 500 + 5;
                                                                    }
int w[maxn][maxn], g[maxn];
                                                                   }
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
                                                                  if (dis[t] == kInf) return false;
w[x][y] += c; w[y][x] += c;
                                                                  for (int i = 0; i < n; i++) h[i] += dis[i];
                                                                  int d = max_flow;
                                                                  for (int u = t; u != s;
pair<int, int> phase(int n) {
memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
int s = -1, t = -1;
                                                                     u = graph[u][p[u]].to) {
                                                                   auto &e = graph[u][p[u]];
                                                                   d = min(d, graph[e.to][e.rev].flow);
while (true) {
 int c = -1;
                                                                 max_flow -= d;
  for (int i = 0; i < n; ++i) {
                                                                  cost += int64_t(d) * h[t];
  if (del[i] || v[i]) continue;
                                                                  for (int u = t; u != s;
```

```
u = graph[u][p[u]].to) {
  auto &e = graph[u][p[u]];
  e.flow += d;
  graph[e.to][e.rev].flow -= d;
}
return true;
}
int MincostMaxflow(
  int ss, int tt, int max_flow, int64_t &cost) {
  this->s = ss, this->t = tt;
  cost = 0;
  fill_n(h, n, 0);
  auto orig_max_flow = max_flow;
  while (Dijkstra(max_flow, cost) && max_flow) {}
  return orig_max_flow - max_flow;
}
};
```

5 Math

5.1 $\lfloor \frac{n}{i} \rfloor$ Enumeration

```
T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor
```

5.2 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b==(x/y)*a;
}
```

5.3 Pollard Rho

```
// does not work when n is prime
// return any non-trivial factor
llu pollard_rho(llu n){
    static auto f=[](llu x,llu k,llu m){
        return add(k,mul(x,x,m),m);
    };
    if (!(n&1)) return 2;
    mt19937 rnd(120821011);
    while(true){
        llu y=2,yy=y,x=rnd()%n,t=1;
        for(llu sz=2;t==1;sz<<=1) {
            for(llu i=0;i<sz;++i){
                if(t!=1)break;
                yy=f(yy,x,n);
                t=gcd(yy>y?yy-y:y-yy,n);
        }
        y=yy;
    }
    if(t!=1&&t!=n) return t;
}
```

5.4 Pi Count (Linear Sieve)

```
static constexpr int N = 1000000 + 5;
11d pi[N];
vector<int> primes;
bool sieved[N];
11d cube_root(11d x){
lld s=cbrt(x-static_cast<long double>(0.1));
while(s*s*s <= x) ++s;
return s-1;
11d square_root(11d x){
lld s=sqrt(x-static_cast<long double>(0.1));
while(s*s <= x) ++s;
return s-1;
void init(){
primes.reserve(N)
primes.push_back(1);
for(int i=2;i<N;i++) {</pre>
 if(!sieved[i]) primes.push_back(i);
  pi[i] = !sieved[i] + pi[i-1];
  for(int p: primes) if(p > 1) {
  if(p * i >= N) break;
   sieved[p * i] = true;
   if(p % i == 0) break;
```

```
11d phi(11d m, 11d n) {
 static constexpr int MM = 80000, NN = 500;
 static lld val[MM][NN];
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
 1ld ret = phi(m,n-1)-phi(m/primes[n],n-1);
 if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
 return ret;
11d pi_count(11d);
11d P2(11d m, 11d n) {
1ld sm = square_root(m), ret = 0;
for(lld i = n+1;primes[i]<=sm;i++)</pre>
  ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
 return ret;
11d pi_count(11d m) {
 if(m < N) return pi[m];</pre>
 11d n = pi_count(cube_root(m));
 return phi(m, n) + n - 1 - P2(m, n);
```

5.5 Strling Number

5.5.1 First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1) \dots (x+n-1) = \sum_{k=0}^n S_1(n,k) x^k$$

$$g(x) = x(x+1) \dots (x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.5.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

5.6 Range Sieve

```
const int MAX_SQRT_B = 50000;
const int MAX_L = 200000 + 5;
bool is_prime_small[MAX_SQRT_B], is_prime[MAX_L];
void sieve(lld 1, lld r){ // [1, r)
for(lld i=2;i*i<r;i++) is_prime_small[i] = true;
for(lld i=1;i<r;i++) is_prime[i-1] = true;
if(l==1) is_prime[0] = false;
for(lld i=2;i*i<r;i++){
   if(!is_prime_small[i]) continue;
   for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;
   for(lld j=std::max(2LL, (l+i-1)/i)*i;j<r;j+=i)
        is_prime[j-1]=false;
}
</pre>
```

5.7 Miller Rabin

```
VL convolution(const VI &a, const VI &b) {
 if(x<2)return 0;</pre>
                                                                  // Should be able to handle N <= 10^5, C <= 10^4
 if(!(x&1))return x==2;
                                                                  int sz = 1;
 llu x1=x-1; int t=0;
                                                                  while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 while(!(x1&1))x1>>=1,t++;
                                                                  vector<cplx> v(sz);
 for(llu m:magic)if(witn(m,x1,x,t))return 0;
                                                                  for (int i = 0; i < sz; ++i) {
                                                                  double re = i < a.size() ? a[i] : 0;</pre>
 return 1;
                                                                  double im = i < b.size() ? b[i] : 0;</pre>
                                                                   v[i] = cplx(re, im);
5.8 Extended Euler
    a^b \equiv \begin{cases} a^b \mod \varphi(m) + \varphi(m) & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases}
                                                                  fft(v, sz);
                                                  (\text{mod } m)
                                                                  for (int i = 0; i <= sz / 2; ++i) {
                                                                  int j = (sz - i) & (sz - 1);
                                                                   cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
5.9 Gauss Elimination
                                                                     * cplx(0, -0.25);
void gauss(vector<vector<double>> &d) {
                                                                   if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
 int n = d.size(), m = d[0].size();
                                                                     ].conj()) * cplx(0, -0.25);
 for (int i = 0; i < m; ++i) {
                                                                   v[i] = x;
  int p = -1;
  for (int j = i; j < n; ++j) {</pre>
                                                                  ifft(v, sz);
   if (fabs(d[j][i]) < eps) continue;</pre>
                                                                  VL c(sz);
   if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
                                                                  for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);
  if (p == -1) continue;
  for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]); for (int j = 0; j < n; ++j) {
                                                                VI convolution_mod(const VI &a, const VI &b, int p) {
                                                                 int sz = 1;
  if (i == j) continue;
                                                                  while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
   double z = d[j][i] / d[i][i];
                                                                  vector<cplx> fa(sz), fb(sz);
   for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
                                                                  for (int i = 0; i < (int)a.size(); ++i)</pre>
                                                                   fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 }
                                                                  for (int i = 0; i < (int)b.size(); ++i)</pre>
                                                                   fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
                                                                  fft(fa, sz), fft(fb, sz);
5.10
      Fast Fourier Transform
                                                                  double r = 0.25 / sz;
                                                                  cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
const int mod = 1000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353
                                                                   int j = (sz - i) & (sz - 1);
                                                                   cplx a1 = (fa[i] + fa[j].conj());
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
                                                                   cplx a2 = (fa[i] - fa[j].conj()) * r2;
  static_assert (M1 <= M2 && M2 <= M3);
                                                                   cplx b1 = (fb[i] + fb[j].conj()) * r3;
 constexpr int64_t r12 = modpow(M1, M2-2, M2);
                                                                   cplx b2 = (fb[i] - fb[j].conj()) * r4;
                                                                  if (i != j) {
  cplx c1 = (fa[j] + fa[i].conj());
 constexpr int64_t r13 = modpow(M1, M3-2, M3);
  constexpr int64_t r23 = modpow(M2, M3-2, M3)
 constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
                                                                    cplx c2 = (fa[j] - fa[i].conj()) * r2;
 B = (B - A + M2) * r12 % M2;

C = (C - A + M3) * r13 % M3;
                                                                    cplx d1 = (fb[j] + fb[i].conj()) * r3;
                                                                    cplx d2 = (fb[j] - fb[i].conj()) * r4;
 C = (C - B + M3) * r23 % M3;
                                                                    fa[i] = c1 * d1 + c2 * d2 * r5;
  return (A + B * M1 + C * M1M2) % mod;
                                                                    fb[i] = c1 * d2 + c2 * d1:
                                                                  fa[j] = a1 * b1 + a2 * b2 * r5;
                                                                  fb[j] = a1 * b2 + a2 * b1;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
                                                                  fft(fa, sz), fft(fb, sz);
const double pi = acos(-1);
                                                                  vector<int> res(sz);
                                                                  for (int i = 0; i < sz; ++i) {
cplx omega[maxn + 1];
                                                                  long long a = round(fa[i].re), b = round(fb[i].re),
void prefft() -
                                                                        c = round(fa[i].im);
 for (int i = 0; i <= maxn; i++)</pre>
                                                                   res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
 omega[i] = cplx(cos(2 * pi * j / maxn),
                                                                 }
     sin(2 * pi * j / maxn));
                                                                  return res;
void fft(vector<cplx> &v, int n) {
                                                                }}
 int z = __builtin_ctz(n) - 1;
                                                                 5.11 Chinese Remainder
 for (int i = 0; i < n; ++i) {
  int x = 0, j = 0;
                                                                lld crt(lld ans[], lld pri[], int n){
  for (;(1 << j) < n;++j) x^{=(i >> j & 1)<<(z - j);
                                                                  lld M = 1, ret = 0;
  if (x > i) swap(v[x], v[i]);
                                                                  for(int i=0;i<n;i++) M *= pri[i];</pre>
                                                                  for(int i=0;i<n;i++){</pre>
                                                                  1ld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
 for (int s = 2; s <= n; s <<= 1) {
  int z = s >> 1;
                                                                  ret += (ans[i]*(M/pri[i])%M * iv)%M;
  for (int i = 0; i < n; i += s) {
                                                                  ret %= M;
   for (int k = 0; k < z; ++k) {
  cplx x = v[i + z + k] * omega[maxn / s * k];
                                                                  return ret;
    v[i + z + k] = v[i + k] - x;
    v[i + k] = v[i + k] + x;
                                                                /*
                                                                Another:
                                                                x = a1 \% m1
                                                                x = a2 \% m2
                                                                g = gcd(m1, m2)
void ifft(vector<cplx> &v, int n) {
                                                                assert((a1-a2)%g==0)
fft(v, n); reverse(v.begin() + 1, v.end());
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
                                                                [p, q] = exgcd(m2/g, m1/g)
                                                                return a2+m2*(p*(a1-a2)/g)
```

0 <= x < lcm(m1, m2)

```
*/
 5.12 Berlekamp Massey
 template <typename T>
 vector<T> BerlekampMassey(const vector<T> &output) {
 vector<T> d(output.size() + 1), me, he;
 for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
   for (size_t j = 0; j < me.size(); ++j)</pre>
   d[i] += output[i - j - 2] * me[j];
if ((d[i] -= output[i - 1]) == 0) continue;
   if (me.empty()) {
    me.resize(f = i);
    continue:
                                                                    return *this;
  vector<T> o(i - f - 1);
T k = -d[i] / d[f]; o.push_back(-k);
                                                                   Poly& imul(LL k) {
  for (T x : he) o.push_back(x * k);
                                                                    return *this:
   if (o.size() < me.size()) o.resize(me.size());</pre>
   for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
  if (i-f+he.size() >= me.size()) he = me, f = i;
  me = o:
 return me;
 5.13 NTT
 template <int mod, int G, int maxn>
 struct NTT {
 static_assert (maxn == (maxn & -maxn));
  int roots[maxn];
 NTT () {
                                                                    Poly Y(*this, _n);
   int r = modpow(G, (mod - 1) / maxn);
   for (int i = maxn >> 1; i; i >>= 1) {
    roots[i] = 1;
    for (int j = 1; j < i; j++)
     roots[i + j] = modmul(roots[i + j - 1], r);
    r = modmul(r, r);
                                                                    return Xi.isz(n());
  // n must be 2^k, and 0 \ll F[i] \ll mod
 void inplace_ntt(int n, int F[], bool inv = false) {
  for (int i = 0, j = 0; i < n; i++) {
  if (i < j) swap(F[i], F[j]);</pre>
    for (int k = n > 1; (j^* = k) < k; k > = 1);
   for (int s = 1; s < n; s *= 2) {
    for (int i = 0; i < n; i += s * 2) {
     for (int j = 0; j < s; j++) {
      int a = F[i+j];
      int b = modmul(F[i+j+s], roots[s+j]);
      F[i+j] = modadd(a, b); // a + b
      F[i+j+s] = modsub(a, b); // a - b
   if (inv) {
                                                                   Poly Dx() const
    int invn = modinv(n);
                                                                   Poly ret(n() - 1);
    for (int i = 0; i < n; i++)
F[i] = modmul(F[i], invn);</pre>
    reverse(F + 1, F + n);
                                                                   Poly Sx() const {
                                                                    Poly ret(n() + 1);
 const int P=2013265921, root=31;
                                                                    return ret;
 const int MAXN=1<<20;</pre>
NTT<P, root, MAXN> ntt;
 5.14 Polynomial Operations
 using VL = vector<LL>;
#define fi(s, n) for (int i=int(s); i<int(n); ++i)
#define Fi(s, n) for (int i=int(n); i>int(s); --i)
int n2k(int n) {
                                                                    if (!_n) return {};
 int sz = 1; while (sz < n) sz <<= 1;</pre>
 return sz;
 template<int MAXN, LL P, LL RT> // MAXN = 2^k
                                                                        tmul(_n, *this)
 struct Poly { // coefficients in [0, P)
 static NTT<MAXN, P, RT> ntt;
 VL coef;
 int n() const { return coef.size(); } // n()>=1
 LL *data() { return coef.data(); }
```

```
const LL *data() const { return coef.data(); }
LL &operator[](size_t i) { return coef[i]; }
const LL &operator[](size_t i)const{return coef[i];}
Poly(initializer_list<LL> a) : coef(a) { }
explicit Poly(int _n = 1) : coef(_n) { }
Poly(const LL *arr, int _n) : coef(arr, arr + _n) {}
Poly(const Poly &p, int _n) : coef(_n) {
 copy_n(p.data(), min(p.n(), _n), data());
Poly& irev(){return reverse(data(),data()+n()),*this;}
Poly& isz(int _n) { return coef.resize(_n), *this; }
Poly& iadd(const Poly &rhs) { // n() == rhs.n()
fi(0, n()) if ((coef[i]+=rhs[i]) >= P)coef[i]-=P;
 fi(0, n()) coef[i] = coef[i] * k % P;
Poly Mul(const Poly &rhs) const {
 const int _n = n2k(n() + rhs.n() - 1);
Poly X(*this, _n), Y(rhs, _n);
ntt(X.data(), _n), ntt(Y.data(), _n);
fi(0, _n) X[i] = X[i] * Y[i] % P;
 ntt(X.data(), _n, true);
return X.isz(n() + rhs.n() - 1);
Poly Inv() const { // coef[0] != 0
if (n() == 1) return {ntt.minv(coef[0])};
const int _n = n2k(n() * 2);
Poly Xi = Poly(*this, (n() + 1)/2).Inv().isz(_n);
 ntt(Xi.data(), _n), ntt(Y.data(), _n);
 fi(0, _n) {
Xi[i] *= (2 - Xi[i] * Y[i]) % P;
  if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
 ntt(Xi.data(), _n, true);
Poly Sqrt() const { // Jacobi(coef[0], P) = 1
 if (n()==1) return {QuadraticResidue(coef[0], P)};
 Poly X = Poly(*this, (n()+1) / 2).Sqrt().isz(n());
 return X.iadd(Mul(X.Inv()).isz(n())).imul(P/2+1);
pair<Poly, Poly> DivMod(const Poly &rhs) const {
 // (rhs.)back() != 0
 if (n() < rhs.n()) return {{0}, *this};</pre>
 const int _n = n() - rhs.n() + 1;
 Poly X(rhs); X.irev().isz(_n);
 Poly Y(*this); Y.irev().isz(_n);
 Poly Q = Y.Mul(X.Inv()).isz(_n).irev();
 X = rhs.Mul(Q), Y = *this;
fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;</pre>
 return {Q, Y.isz(max(1, rhs.n() - 1))};
 fi(0, ret.n()) ret[i] = (i + 1) * coef[i + 1] % P;
 return ret.isz(max(1, ret.n()));
 fi(0, n()) ret[i + 1]=ntt.minv(i + 1)*coef[i] % P;
Poly _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
 return Poly(Y.data() + n() - 1, nn);
VL _eval(const VL &x, const auto up)const{
const int _n = (int)x.size();
 vector<Poly> down(_n * 2);
 down[1] = DivMod(up[1]).second;
 fi(2,_n*2) down[i]=down[i/2].DivMod(up[i]).second;
 /* down[1] = Poly(up[1]).irev().isz(n()).Inv().irev()
 fi(2, _n * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
   1, down[i / 2]); */
 VL y(_n);
 fi(0, _n) y[i] = down[_n + i][0];
```

```
*z = (x0y0 + x1y1 , x0y1 + x1y0 )
  return y:
                                                             * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
static vector<Poly> _tree1(const VL &x) {
                                                             * z' = ((x\theta+x1)(y\theta+y1)), (x\theta-x1)(y\theta-y1))
* z = (1/2) * z''
 const int _n = (int)x.size();
  vector<Poly> up(_n * 2);
                                                             * or convolution:
  fi(0, _n) up[_n + i] = \{(x[i] ? P - x[i] : 0), 1\};
 Fi(0, _n-1) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
                                                             * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                             * and convolution:
  return up:
                                                             * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
VL Eval(const VL&x)const{return _eval(x,_tree1(x));}
                                                            const LL MOD = 1e9+7;
static Poly Interpolate(const VL &x, const VL &y) {
                                                            inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
 const int _n = (int)x.size();
                                                             for( int d = 1 ; d < N ; d <<= 1 ) {</pre>
  vector<Poly> up = _{tree1}(x), _{down(_n * 2)};
                                                              int d2 = d << 1;
 VL z = up[1].Dx()._eval(x, up);
                                                              for( int s = 0 ; s < N ; s += d2 )
 fi(0, _n) z[i] = y[i] * ntt.minv(z[i]) % P;
                                                               x[ i ] = ta+tb;
                                                                x[ j ] = ta-tb;
   .iadd(down[i * 2 + 1].Mul(up[i * 2]));
                                                                if( x[ i ] >= MOD ) x[ i ] -= MOD;
  return down[1];
                                                                if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
Poly Ln() const { // coef[0] == 1
  return Dx().Mul(Inv()).Sx().isz(n());
                                                             if( inv )
                                                              for( int i = 0 ; i < N ; i++ ) {</pre>
Poly Exp() const \{ // coef[0] == 0 \}
 if (n() == 1) return {1};
Poly X = Poly(*this, (n() + 1)/2).Exp().isz(n());
Poly Y = X.Ln(); Y[0] = P - 1;
                                                               x[ i ] *= inv( N, MOD );
                                                               x[ i ] %= MOD;
  fi(0, n()) if((Y[i] = coef[i] - Y[i]) < 0)Y[i]+=P;
 return X.Mul(Y).isz(n());
                                                                  DiscreteLog
                                                            5.16
Poly Pow(const string &K) const {
                                                            template<typename Int>
                                                            Int BSGS(Int x, Int y, Int M) {
 int nz = 0:
  while (nz < n() && !coef[nz]) ++nz;</pre>
                                                              // x^? \equiv y (mod M)
                                                              Int t = 1, c = 0, g = 1;
 LL nk = 0, nk2 = 0;
                                                              for (Int M_ = M; M_ > 0; M_ >>= 1)
  for (char c : K) {
  nk = (nk * 10 + c - '0') % P;
                                                                g = g * x % M;
  nk2 = nk2 * 10 + c - '0';
                                                              for (g = gcd(g, M); t % g != 0; ++c) {
   if (nk2 * nz >= n()) return Poly(n());
                                                                if (t == y) return c;
   nk2 %= P - 1;
                                                                t = t * x % M;
  if (!nk && !nk2) return Poly({1}, n());
                                                              if (y % g != 0) return -1;
                                                              t /= g, y /= g, M /= g;
Int h = 0, gs = 1;
 Poly X(data() + nz, n() - nz * nk2);
 LL \times 0 = X[0]
                                                              for (; h * h < M; ++h) gs = gs * x % M;
  return X.imul(ntt.minv(x0)).Ln().imul(nk).Exp()
   .imul(ntt.mpow(x0, nk2)).irev().isz(n()).irev();
                                                              unordered_map<Int, Int> bs;
                                                              for (Int s = 0; s < h; bs[y] = ++s)
Poly InvMod(int L) { // (to evaluate linear recursion)
                                                                y = y * x % M;
 Poly R{1, 0}; // *this * R mod x^L = 1 (*this[0] ==
                                                              for (Int s = 0; s < M; s += h) {
                                                                t = t * gs % M;
  for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                                if (bs.count(t)) return c + s + h - bs[t];
  Poly 0 = R.Mul(Poly(data(), min(2 \ll level, n())));
   Poly Q(2 << level); Q[0] = 1;
                                                              return -1;
  for (int j = (1 << level); j < (2 << level); ++j)
Q[j] = (P - O[j]) % P;</pre>
                                                            5.17
                                                                   FloorSum
   R = R.Mul(Q).isz(4 << level);
                                                            // @param n `n < 2^32`
                                                            // @param m `1 <= m < 2^32`
  return R.isz(L);
                                                            // @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
static LL LinearRecursion(const VL&a,const VL&c,LL n){
                                                            llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
 // a_n = \sum_{j=0}^{n} a_{n-j}
                                                             11u ans = 0:
 const int k = (int)a.size();
                                                             while (true)
  assert((int)c.size() == k + 1);
                                                              if (a >= m) {
 Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
                                                               ans += n * (n - 1) / 2 * (a / m); a %= m;
  fi(1, k + 1) C[k - i] = c[i] ? P - c[i] : 0;
 C[k] = 1
                                                              if (b >= m) {
 while (n) {
                                                               ans += n * (b / m); b %= m;
  if (n % 2) W = W.Mul(M).DivMod(C).second;
   n /= 2, M = M.Mul(M).DivMod(C).second;
                                                              llu y_max = a * n + b;
                                                              if (y_max < m) break;</pre>
 LL ret = 0:
                                                              // y_max < m * (n + 1)
  fi(0, k) ret = (ret + W[i] * a[i]) % P;
                                                              // floor(y_max / m) <= n
 return ret;
                                                              n = (11u)(y_max / m), b = (11u)(y_max % m);
}
                                                              swap(m, a);
#undef fi
                                                             return ans;
#undef Fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
                                                            lld floor_sum(lld n, lld m, lld a, lld b) {
template<> decltype(Poly_t::ntt) Poly_t::ntt = {};
                                                             llu ans = 0;
                                                             if (a < 0) {
5.15 FWT
                                                              11u \ a2 = (a \% m + m) \% m;
/* xor convolution:
                                                              ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
* x = (x0, x1) , y = (y0, y1)
                                                              a = a2;
```

```
if (b < 0) {
 11\dot{u} b2 = (b % m + m) % m;
 ans -= 1ULL * n * ((b2 - b) / m);
b = b2
return ans + floor_sum_unsigned(n, m, a, b);
```

5.18 ExtendedFloorSum

```
g(a,b,c,n) = \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor
                              a \geq c \vee b \geq c
                                                                                                                     n < 0 \lor a = 0
                               \begin{bmatrix} \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), \end{bmatrix} 
                                                                                                                       otherwise
h(a,b,c,n) = \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2
                              \left( \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right)
                                +\lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1)
                                +h(a\bmod c,b\bmod c,c,n)
                                +2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)
                             \left\{+2\left\lfloor\frac{b}{c}\right\rfloor\cdot f(a\ \mathsf{mod}\ c,b\ \mathsf{mod}\ c,c,n),\right.
                                                                                                                      a \geq c \vee b \geq c
```

nm(m+1) - 2g(c, c-b-1, a, m-1)

-2f(c, c-b-1, a, m-1) - f(a, b, c, n), otherwise

5.19 Quadratic residue

```
struct S {
 int MOD, w;
 int64_t x, y;
 S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
 : MOD(m), w(w_), x(x_), y(y_) {}
S operator*(const S &rhs) const {
  int w_{-} = w;
  if (w_ == -1) w_ = rhs.w;
  assert(w_ != -1 and w_ == rhs.w);
  return { MOD, w_,
(x * rhs.x + y * rhs.y % MOD * w) % MOD,
   (x * rhs.y + y * rhs.x) % MOD };
};
int get_root(int n, int P) {
  if (P == 2 or n == 0) return n;
  if (qpow(n, (P - 1) / 2, P) != 1) return -1;
  auto check = [&](int x) {
  return qpow(x, (P - 1) / 2, P); };
if (check(n) == P-1) return -1;
  int64_t a; int w; mt19937 rnd(7122);
  do { a = rnd() % P;
w = ((a * a - n) % P + P) % P;
  } while (check(w) != P - 1);
  return qpow(S(P, w, a, 1), (P + 1) / 2).x;
```

5.20 De-Bruijn

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
  if (n \% p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
 } else {
 aux[t] = aux[t - p];
db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {
   aux[t] = i;
   db(t + 1, t, n, k);
int de_bruijn(int k, int n) {
// return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
if (k == 1) {
 res[0] = 0;
  return 1;
```

```
for (int i = 0; i < k * n; i++) aux[i] = 0;
db(1, 1, n, k);
return sz;
```

5.21 Simplex Construction

Standard form: maximize $\sum_{1 \leq i \leq n} c_i x_i$ such that for all $1 \leq j \leq m$, $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$ and $x_i \ge 0$ for all $1 \le i \le n$.

- 1. In case of minimization, let $c_i^\prime = -c_i$
- 2. $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j$
- $3. \sum_{1 < i < n} A_{ji} x_i = b_j$
 - $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$
 - $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

5.22 Simplex

 $n < 0 \lor a = 0$

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return VD(n, -inf) if the solution doesn't exist // return VD(n, +inf) if the solution is unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
 double inv = 1.0 / d[r][s];
 for (int i = 0; i < m + 2; ++i)
  for (int j = 0; j < n + 2; ++j)
   if (i != r && j != s)
    d[i][j] = d[r][j] * d[i][s] * inv;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;
 d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
 int x = m + z;
 while (true) {
  int s = -1;
  for (int i = 0; i <= n; ++i) {
  if (!z && q[i] == -1) continue;
   if (s == -1) \mid d[x][i] < d[x][s]) s = i;
  if (d[x][s] > -eps) return true;
  int r = -1;
  for (int i = 0; i < m; ++i) {</pre>
   if (d[i][s] < eps) continue;
if (r == -1 || \</pre>
     d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
  pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
 m = b.size(), n = c.size();
 d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)</pre>
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)
  p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
 for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i]; q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
 for (int i = 1; i < m; ++i)
if (d[i][n + 1] < d[r][n + 1]) r = i;
 if (d[r][n + 1] < -eps) {</pre>
  pivot(r, n);
  if (!phase(1) \mid | d[m + 1][n + 1] < -eps)
   return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
   int s = min_element(d[i].begin(), d[i].end() - 1)
        - d[i].begin();
```

```
pivot(i, s);
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)
  if (p[i] < n) x[p[i]] = d[i][n + 1];
 return x:
5.23 Charateristic Polynomial
vector<vector<int>> Hessenberg(const vector<vector<int
     >> &A) {
 int N = A.size();
 vector<vector<int>> H = A;
 for (int i = 0; i < N - 2; ++i) {
  if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {
    if (H[j][i]) {
     for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
      for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k
     ][j]);
     break:
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
   * H[i + 1][k] * (kP - coef)) % kP;
for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
  }
 return H;
vector<int> CharacteristicPoly(const vector<vector<int
    >> &A) {
 int N = A.size();
 auto H = Hessenberg(A);
for (int i = 0; i < N; ++i) {</pre>
  for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
 vector<vector<int>>> P(N + 1, vector<int>(N + 1));
 P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {</pre>
  P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j -
    1];
  int val = 1;
  for (int j = i - 1; j >= 0; --j) {
  int coef = 1LL * val * H[j][i - 1] % kP;
   for (int k = 0; k \le j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 if (N & 1) {
  for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
 return P[N];
5.24 Partition Number
int b = sqrt(n);
ans[0] = tmp[0] = 1;
for (int i = 1; i <= b; i++) {
for (int rep = 0; rep < 2; rep++)
for (int j = i; j <= n - i * i; j++)
modadd(tmp[j], tmp[j-i]);
for (int j = i * i; j <= n; j++)</pre>
  modadd(ans[j], tmp[j - i * i]);
```

6 Geometry

6.1 Basic Geometry

```
using coord_t = int;
using Real = double;
using Point = std::complex<coord_t>;
int sgn(coord_t x) {
return (x > 0) - (x < 0); }
coord_t dot(Point a, Point b) {
 return real(conj(a) * b); }
coord_t cross(Point a, Point b) {
 return imag(conj(a) * b); }
int ori(Point a, Point b, Point c) {
 return sgn(cross(b - a, c - a)); }
bool operator<(const Point &a, const Point &b) {</pre>
 return real(a) != real(b)
  ? real(a) < real(b) : imag(a) < imag(b);
int argCmp(Point a, Point b) {
  // -1 / 0 / 1 <-> < / == / > (atan2)
 int qa = (imag(a) == 0
   ? (real(a) < 0 ? 3 : 1) : (imag(a) < 0 ? 0 : 2));
 int qb = (imag(b) == 0
   ? (real(b) < 0 ? 3 : 1) : (imag(b) < 0 ? 0 : 2));
 if (qa != qb) return sgn(qa - qb);
 return sgn(cross(b, a));
template <typename V> Real area(const V & pt) {
 coord_t ret = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 return ret / 2.0;
6.2 2D Convex Hull
template<typename PT>
vector<PT> buildConvexHull(vector<PT> d) {
 sort(ALL(d), [](const PT& a, const PT& b){
   return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
 vector<PT> s(SZ(d)<<1);</pre>
 int o = 0;
 for(auto p: d) {
  while(o>=2 && cross(p-s[o-2], s[o-1]-s[o-2])<=0)
  s[o++] = p;
 for(int i=SZ(d)-2, t = o+1; i>=0; i--){
  while(o>=t&&cross(d[i]-s[o-2],s[o-1]-s[o-2])<=0)
  s[o++] = d[i];
 s.resize(o-1);
 return s;
6.3
     3D Convex Hull
// return the faces with pt indexes
int flag[MXN][MXN];
struct Point{
 ld x,y,z;
 Point operator * (const 1d &b) const {
  return (Point) {x*b,y*b,z*b};}
 Point operator * (const Point &b) const {
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
Point ver(Point a, Point b, Point c) {
  return (b - a) * (c - a);}
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0;
 REP(i,n) REP(j,n) flag[i][j] = 0;
 vector<Face> now
 now.emplace_back(0,1,2);
 now.emplace_back(2,1,0);
for (int i=3; i<n; i++){
  ftop++; vector<Face> next;
  REP(j, SZ(now)) {
   Face& f=now[j]; int ff = 0;
   ld d=(pt[i]-pt[f.a]).dot(
     ver(pt[f.a], pt[f.b], pt[f.c]));
   if (d <= 0) next.push_back(f);</pre>
   if (d > 0) ff=ftop;
   else if (d < 0) ff=-ftop;
flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;</pre>
```

```
REP(j, SZ(now)) {
                                                                   if (dis(v[p], v[i]) < d) {</pre>
   Face& f=now[j]
                                                                    d = dis(v[p], v[i]);
   if (flag[f.a][f.b] > 0 &&
                                                                    found = true;
     flag[f.a][f.b] != flag[f.b][f.a])
    next.emplace_back(f.a,f.b,i);
   if (flag[f.b][f.c] > 0 &&
                                                                 }
     flag[f.b][f.c] != flag[f.c][f.b])
    next.emplace_back(f.b,f.c,i);
                                                               if (found) rebuild_m(i + 1);
   if (flag[f.c][f.a] > 0 &&
                                                                else m[kx][ky][kz] = i;
     flag[f.c][f.a] != flag[f.a][f.c])
    next.emplace_back(f.c,f.a,i);
                                                               return d;
                                                             }
 now=next:
                                                             6.7
                                                                  Simulated Annealing
 return now;
                                                             11f anneal() {
                                                              mt19937 rnd_engine( seed );
     2D Farthest Pair
                                                               uniform_real_distribution< llf > rnd( 0, 1 );
6.4
                                                               const llf dT = 0.001;
// stk is from convex hull
                                                               // Argument p
n = (int)(stk.size());
                                                               1lf S_cur = calc( p ), S_best = S_cur;
for ( 1lf T = 2000 ; T > EPS ; T -= dT ) {
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {</pre>
                                                                // Modify p to p_prime
 while(abs(cross(stk[i+1]-stk[i],
                                                                const llf S_prime = calc( p_prime );
   stk[(pos+1)%n]-stk[i])) >
                                                               const llf delta_c = S_prime - S_cur;
llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
   abs(cross(stk[i+1]-stk[i],
   stk[pos]-stk[i]))) pos = (pos+1)%n;
                                                               if ( rnd( rnd_engine ) <= prob )</pre>
 ans = max({ans, dis(stk[i], stk[pos]),
                                                                S_cur = S_prime, p = p_prime;
  dis(stk[i+1], stk[pos])});
                                                                if ( S_prime < S_best ) // find min</pre>
                                                                 S_best = S_prime, p_best = p_prime;
6.5 2D Closest Pair
                                                               return S_best;
struct cmp_y {
 bool operator()(const P& p, const P& q) const {
  return p.y < q.y;</pre>
                                                              6.8 Half Plane Intersection
                                                             // NOTE: Point is complex<Real>
multiset<P, cmp_y> s;
                                                              // cross(pt-line.st, line.dir)<=0 <-> pt in half plane
void solve(P a[], int n) {
                                                             struct Line {
 sort(a, a + n, [](const P& p, const P& q) {
                                                               Point st, ed;
  return tie(p.x, p.y) < tie(q.x, q.y);</pre>
                                                                Point dir;
                                                               Line (Point _s, Point _e)
 11f d = INF; int pt = 0;
                                                                 : st(_s), ed(_e), dir(_e - _s) {}
 for (int i = 0; i < n; ++i) {
 while (pt < i and a[i].x - a[pt].x >= d)
                                                             bool operator<(const Line &lhs, const Line &rhs) {</pre>
   s.erase(s.find(a[pt++]));
                                                                if (int cmp = argCmp(lhs.dir, rhs.dir))
  auto it = s.lower_bound(P(a[i].x, a[i].y - d));
                                                                  return cmp == -1;
  while (it != s.end() and it->y - a[i].y < d)
                                                                return ori(lhs.st, lhs.ed, rhs.st) < 0;</pre>
   d = min(d, dis(*(it++), a[i]));
  s.insert(a[i]);
                                                             Point intersect(const Line &A, const Line &B) {
                                                               Real t = cross(B.st - A.st, B.dir) /
                                                                cross(A.dir, B.dir)
                                                                return A.st + t * A.dir;
6.6 kD Closest Pair (3D ver.)
                                                             Real HPI(vector<Line> &lines) {
11f solve(vector<P> v) {
                                                                sort(lines.begin(), lines.end());
 shuffle(v.begin(), v.end(), mt19937());
                                                                deque<Line> que;
 unordered_map<lld, unordered_map<lld,
                                                                deque<Point> pt;
  unordered_map<lld, int>>> m;
                                                                que.push_back(lines[0]);
 llf d = dis(v[0], v[1]);
                                                                for (int i = 1; i < (int)lines.size(); i++) {</pre>
 auto Idx = [&d] (11f x) -> 11d {
                                                                  if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
  return round(x * 2 / d) + 0.1;
                                                                   continue;
 auto rebuild_m = [&m, &v, &Idx](int k) {
                                                             #define POP(L, R) \
  m.clear():
                                                                  while (pt.size() > 0 \
  for (int i = 0; i < k; ++i)
m[Idx(v[i].x)][Idx(v[i].y)]</pre>
                                                                    && ori(L.st, L.ed, pt.back()) < 0) \
                                                                  pt.pop_back(), que.pop_back(); \
while (pt.size() > 0 \
    [Idx(v[i].z)] = i;
 }; rebuild_m(2);
                                                                    && ori(R.st, R.ed, pt.front()) < 0) \
 for (size_t i = 2; i < v.size(); ++i) {</pre>
                                                                    pt.pop_front(), que.pop_front();
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
                                                                  POP(lines[i], lines[i]);
     kz = Idx(v[i].z); bool found = false;
                                                                  pt.push_back(intersect(que.back(), lines[i]));
  for (int dx = -2; dx <= 2; ++dx) {
                                                                  que.push_back(lines[i]);
   const 11d nx = dx + kx
   if (m.find(nx) == m.end()) continue;
                                                               POP(que.front(), que.back())
   auto& mm = m[nx];
                                                                if (que.size() <= 1 ||</pre>
   for (int dy = -2; dy <= 2; ++dy) {
                                                                  argCmp(que.front().dir, que.back().dir) == 0)
    const 11d ny = dy + ky;
                                                                  return 0:
    if (mm.find(ny) == mm.end()) continue;
                                                               pt.push_back(intersect(que.front(), que.back()));
    auto& mmm = mm[ny];
                                                                return area(pt);
    for (int dz = -2; dz <= 2; ++dz) {
     const 11d nz = dz + kz;
     if (mmm.find(nz) == mmm.end()) continue;
                                                                   Minkowski Sum
     const int p = mmm[nz];
```

// sign1 = 1 for outer tang, -1 for inter tang

```
vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
                                                               vector<Line> ret;
hull(A), hull(B);
                                                               double d_{sq} = norm2(c1.0 - c2.0);
 vector<pll> C(1, A[0] + B[0]), s1, s2;
                                                               if( d_sq < eps ) return ret;</pre>
 for(int i = 0; i < SZ(A); ++i)
s1.pb(A[(i + 1) % SZ(A)] - A[i]);
                                                               double d = sqrt( d_sq )
                                                               Pt v = (c2.0 - c1.0) / d;
 for(int i = 0; i < SZ(B); i++)
s2.pb(B[(i + 1) % SZ(B)] - B[i]);
                                                               double c = (c1.R - sign1 * c2.R) / d;
                                                               if( c * c > 1 ) return ret;
 for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
                                                               double h = sqrt( max( 0.0 , 1.0 - c * c ) );
                                                               for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
  Pt n = { v.X * c - sign2 * h * v.Y ,
  if (p2 >= SZ(B)
    || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
                                                                 v.Y * c + sign2 * h * v.X };
   C.pb(C.back() + s1[p1++]);
                                                                Pt p1 = c1.0 + n * c1.R;
                                                                Pt p2 = c2.0 + n * (c2.R * sign1);
   C.pb(C.back() + s2[p2++]);
 return hull(C), C;
                                                                if( fabs( p1.X - p2.X ) < eps and
                                                                  fabs( p1.Y - p2.Y ) < eps )
                                                                 p2 = p1 + perp(c2.0 - c1.0);
6.10 Intersection of line and Circle
                                                                ret.push_back( { p1 , p2 } );
vector<pdd> line_interCircle(const pdd &p1,
                                                               return ret;
    const pdd &p2,const pdd &c,const double r){
                                                              }
 pdd ft=foot(p1,p2,c),vec=p2-p1;
 double dis=abs(c-ft);
                                                                     Minimum Covering Circle
 if(fabs(dis-r)<eps) return vector<pdd>{ft};
                                                              template<typename P>
 if(dis>r) return {};
                                                              Circle getCircum(const P &a, const P &b, const P &c){
 vec=vec*sqrt(r*r-dis*dis)/abs(vec);
                                                               Real a1 = a.x-b.x, b1 = a.y-b.y;
 return vector<pdd>{ft+vec,ft-vec};
                                                               Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
                                                               Real a2 = a.x-c.x, b2 = a.y-c.y;
                                                               Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
6.11 Intersection of Polygon and Circle
                                                               Circle cc;
// Divides into multiple triangle, and sum up
                                                               cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
// test by HDU2892
                                                               cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
const double PI=acos(-1);
                                                               cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
                                                               return cc:
 if(abs(pb)<eps) return 0;</pre>
 double S, h, theta;
                                                              template<typename P>
 double a=abs(pb), b=abs(pa), c=abs(pb-pa);
                                                              Circle MinCircleCover(const vector<P>& pts){
 double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
                                                               random_shuffle(pts.begin(), pts.end());
 double cosC = dot(pa,pb) / a / b, C = acos(cosC);
                                                               Circle c = { pts[0], 0 };
 if(a > r){
                                                               for(int i=0;i<(int)pts.size();i++){</pre>
 S = (C/2)*r*r;
                                                                if (dist(pts[i], c.o) <= c.r) continue;</pre>
 h = a*b*sin(C)/c;
                                                                c = { pts[i], 0 };
for (int j = 0; j < i; j++) {
  if (h < r && B < PI/2)
   S = (acos(h/r)*r*r - h*sqrt(r*r-h*h));
                                                                 if(dist(pts[j], c.o) <= c.r) continue;</pre>
                                                                 c.o = (pts[i] + pts[j]) / 2;
 else if(b > r){
                                                                 c.r = dist(pts[i], c.o)
  theta = PI - B - asin(sin(B)/r*a);
                                                                 for (int k = 0; k < j; k++) {
  S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
                                                                  if (dist(pts[k], c.o) <= c.r) continue;</pre>
                                                                  c = getCircum(pts[i], pts[j], pts[k]);
 else S = .5*sin(C)*a*b;
 return S;
                                                                }
double area_poly_circle(const vector<pdd> poly,
                                                               return c;
  const pdd &0,const double r){
 double S=0;
 for(int i=0;i<SZ(poly);++i)</pre>
                                                                    KDTree (Nearest Point)
                                                              6.15
  S+=_area(poly[i]-0,poly[(i+1)%SZ(poly)]-0,r)
                                                              const int MXN = 100005;
    *ori(0,poly[i],poly[(i+1)%SZ(poly)]);
                                                              struct KDTree {
 return fabs(S);
                                                               struct Node {
}
                                                                int x,y,x1,y1,x2,y2;
                                                                int id,f;
Node *L, *R;
6.12 Intersection of Two Circle
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
                                                               } tree[MXN], *root;
 pdd o1 = a.0, o2 = b.0;
                                                               int n;
 double r1 = a.R, r2 = b.R, d2 = abs2(o1 - o2),
                                                               LL dis2(int x1, int y1, int x2, int y2) {
     d = sqrt(d2)
                                                                LL dx = x1-x2, dy = y1-y2;
 if(d < max(r1, r2) - min(r1, r2) \mid | d > r1 + r2)
                                                                return dx*dx+dy*dy;
  return 0;
 pdd u = (o1 + o2) * 0.5
                                                               static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
  + (o1 - o2) * ((r2 * r2 - r1 * r1) / (2 * d2));
                                                               static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 double A = sqrt((r1 + r2 + d) * (r1 - r2 + d)
                                                               void init(vector<pair<int,int>> ip) {
     * (r1 + r2 - d) * (-r1 + r2 + d));
                                                                n = ip.size();
 pdd v = pdd(o1.Y - o2.Y, -o1.X + o2.X) * A
                                                                for (int i=0; i<n; i++) {</pre>
  / (2 * d2);
                                                                 tree[i].id = i;
 p1 = u + v, p2 = u - v;
                                                                 tree[i].x = ip[i].first;
 return 1;
                                                                 tree[i].y = ip[i].second;
                                                                root = build_tree(0, n-1, 0);
     Tangent line of Two Circle
6.13
vector<Line> go(const Cir& c1,
                                                               Node* build_tree(int L, int R, int d) {
  const Cir& c2, int sign1){
                                                                if (L>R) return nullptr;
                                                                int M = (L+R)/2; tree[M].f = d%2;
```

```
nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
                                                                 int C;
  tree[M].x1 = tree[M].x2 = tree[M].x;
                                                                 Cir c[N]
                                                                 bool g[N][N], overlap[N][N];
  tree[M].y1 = tree[M].y2 = tree[M].y;
                                                                 // Area[i] : area covered by at least i circles
double Area[ N ];
  tree[M].L = build_tree(L, M-1, d+1);
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
                                                                 void init(int _C){ C = _C;}
   tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
                                                                 struct Teve {
                                                                  pdd p; double ang; int add;
   tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
                                                                  Teve() {}
                                                                  Teve(pdd _a, double _b, int _c):p(_a), ang(_b), add(
  tree[M].R = build_tree(M+1, R, d+1);
  if (tree[M].R) {
                                                                  bool operator<(const Teve &a)const
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
                                                                  {return ang < a.ang;}
                                                                 }eve[N * 2];
   tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
                                                                 // strict: x = 0, otherwise x = -1
                                                                 bool disjuct(Cir &a, Cir &b, int x)
   tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
                                                                 {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
                                                                 bool contain(Cir &a, Cir &b, int x)
  return tree+M;
                                                                 \{return sign(a.R - b.R - abs(a.0 - b.0)) > x;\}
 int touch(Node* r, int x, int y, LL d2){
                                                                 bool contain(int i, int j) {
 LL dis = sqrt(d2)+1;
                                                                  /* c[j] is non-strictly in c[i]. */
  if (x<r->x1-dis || x>r->x2+dis ||
                                                                  return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c[i].R) \mid c[i].R - c[i].R)
                                                                     [j].R) == 0 && i < j)) && contain(c[i], c[j], -1);
    y<r->y1-dis || y>r->y2+dis)
   return 0:
                                                                 void solve(){
  return 1:
                                                                  fill_n(Area, C + 2, 0);
                                                                  for(int i = 0; i < C; ++i)
 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
 if (!r || !touch(r, x, y, md2)) return;
LL d2 = dis2(r->x, r->y, x, y);
                                                                   for(int j = 0; j < C; ++j)
                                                                    overlap[i][j] = contain(i, j);
  if (d2 < md2 || (d2 == md2 && mID < r->id)) {
                                                                  for(int i = 0; i < C; ++i)
  mID = r->id;
                                                                   for(int j = 0; j < C; ++j)
                                                                    g[i][j] = !(overlap[i][j] || overlap[j][i] ||
   md2 = d2;
                                                                       disjuct(c[i], c[j], -1));
  // search order depends on split dim
                                                                  for(int i = 0; i < C; ++i){</pre>
  if ((r->f == 0 \&\& x < r->x) ||
                                                                   int E = 0, cnt = 1;
    (r->f == 1 \&\& y < r->y)) {
                                                                   for(int j = 0; j < C; ++j)</pre>
                                                                    if(j != i && overlap[j][i])
   nearest(r->L, x, y, mID, md2);
   nearest(r->R, x, y, mID, md2);
                                                                     ++cnt;
                                                                   for(int j = 0; j < C; ++j)
if(i != j && g[i][j]) {</pre>
  } else {
   nearest(r->R, x, y, mID, md2);
   nearest(r->L, x, y, mID, md2);
                                                                     pdd aa, bb;
                                                                     CCinter(c[i], c[j], aa, bb);
llf A = atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);
 int query(int x, int y) {
                                                                      11f B = atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);
  int id = 1029384756;
                                                                     eve[E++] = Teve(bb,B,1), eve[E++]=Teve(aa,A,-1);
  LL d2 = 102938475612345678LL;
                                                                     if(B > A) ++cnt;
  nearest(root, x, y, id, d2);
                                                                   if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
  return id;
} tree;
                                                                    sort(eve, eve + E);
                                                                    eve[E] = eve[0];
6.16
      Rotating Sweep Line
                                                                    for(int j = 0; j < E; ++j){}
                                                                     cnt += eve[j].add;
void rotatingSweepLine(pair<int, int> a[], int n) {
 vector<pair<int, int>> 1;
                                                                     Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
 1.reserve(n * (n - 1) / 2)
                                                                      double theta = eve[j + 1].ang - eve[j].ang;
 for (int i = 0; i < n; ++i)
                                                                     if (theta < 0) theta += 2. * pi:</pre>
  for (int j = i + 1; j < n; ++j)
                                                                     Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
   1.emplace_back(i, j);
 sort(l.begin(), l.end(), [&a](auto &u, auto &v){
  1ld udx = a[u.first].first - a[u.second].first;
  lld udy = a[u.first].second - a[u.second].second;
lld vdx = a[v.first].first - a[v.second].first;
  11d vdy = a[v.first].second - a[v.second].second;
  if (udx == 0 \text{ or } vdx == 0) \text{ return not } udx == 0;
                                                                     Stringology
  int s = sgn(udx * vdx);
  return udy * vdx * s < vdy * udx * s;
                                                                7.1 Hash
 });
                                                                class Hash {
 vector<int> idx(n), p(n);
 iota(idx.begin(), idx.end(), 0);
                                                                 private:
                                                                  static constexpr int P = 127, Q = 1051762951;
 sort(idx.begin(), idx.end(), [&a](int i, int j){
 return a[i] < a[j]; });
for (int i = 0; i < n; ++i) p[idx[i]] = i;
                                                                  vector<int> h, p;
                                                                 public:
                                                                  void init(const string &s){
 for (auto [i, j]: 1) {
                                                                   h.assign(s.size()+1, 0); p.resize(s.size()+1);
  // do here
                                                                   for (size_t i = 0; i < s.size(); ++i)</pre>
  swap(p[i], p[j]);
                                                                    h[i + 1] = add(mul(h[i], P), s[i]);
  idx[p[i]] = i, idx[p[j]] = j;
                                                                   generate(p.begin(), p.end(),[x=1,y=1,this]()
                                                                     mutable{y=x;x=mul(x,P);return y;});
6.17 Circle Cover
                                                                  int query(int 1, int r){ // 1-base (1, r]
                                                                   return sub(h[r], mul(h[1], p[r-1]));}
const int N = 1021;
struct CircleCover {
```

7.2 Suffix Arrau namespace sfxarray { bool t[maxn * 2]; int hi[maxn], rev[maxn]; int _s[maxn * 2], sa[maxn * 2], c[maxn * 2]; int x[maxn], p[maxn], q[maxn * 2]; // sa[i]: sa[i]-th suffix is the \ // i-th lexigraphically smallest suffix. // hi[i]: longest common prefix \ // of suffix sa[i] and suffix sa[i - 1]. void pre(int *sa, int *c, int n, int z) { memset(sa, 0, sizeof(int) * n); memcpy(x, c, sizeof(int) * z); void induce(int *sa,int *c,int *s,bool *t,int n,int z){ memcpy(x + 1, c, sizeof(int) * (z - 1)); for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] - 1]) sa[x[s[sa[i] - 1]]++] = sa[i] - 1;memcpy(x, c, sizeof(int) * z); for (int i = n - 1; $i \ge 0$; --i) if (sa[i] && t[sa[i] - 1]) sa[--x[s[sa[i] - 1]]] = sa[i] - 1;void sais(int *s, int *sa, int *p, int *q, bool *t, int *c, int n, int z) { bool uniq = t[n - 1] = true; int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1; memset(c, 0, sizeof(int) * z); for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2; for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i]; if (uniq) { for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i; return: for (int i = n - 2; i >= 0; --i) t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);pre(sa, c, n, z); for (int i = 1; i <= n - 1; ++i)</pre> if (t[i] && !t[i - 1]) sa[--x[s[i]]] = p[q[i] = nn++] = i;induce(sa, c, s, t, n, z); for (int i = 0; i < n; ++i) if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) { bool neq = last < 0 ||</pre> memcmp(s + sa[i], s + last,(p[q[sa[i]] + 1] - sa[i]) * sizeof(int)); ns[q[last = sa[i]]] = nmxz += neq; }} sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1); pre(sa, c, n, z); for (int i = nn - 1; i >= 0; --i) sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];induce(sa, c, s, t, n, z); void build(const string &s) { for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre> _s[(int)s.size()] = 0; // s shouldn't contain 0 sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256); for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];</pre> for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre> int ind = 0; hi[0] = 0; for (int i = 0; i < (int)s.size(); ++i) { if (!rev[i]) {</pre> ind = 0; continue: while (i + ind < (int)s.size() && \</pre> s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;hi[rev[i]] = ind ? ind-- : 0; }} 7.3 Suffix Automaton struct SuffixAutomaton { struct node int ch[K], len, fail, cnt, indeg; $node(int L = 0) : ch{}, len(L), fail(0), cnt(0),$

return f; vector < int > f = kmp(t), r

return res;

```
indeg(0) {}
} st[N];
int root, last, tot;
```

```
void extend(int c) {
  int cur = ++tot;
  st[cur] = node(st[last].len + 1);
  while (last && !st[last].ch[c]) {
    st[last].ch[c] = cur;
    last = st[last].fail;
  if (!last) {
    st[cur].fail = root;
  } else {
    int q = st[last].ch[c];
    if (st[q].len == st[last].len + 1) {
      st[cur].fail = q;
    } else {
      int clone = ++tot;
      st[clone] = st[q];
st[clone].len = st[last].len + 1;
      st[st[cur].fail = st[q].fail = clone].cnt = 0;
      while (last && st[last].ch[c] == q) {
        st[last].ch[c] = clone;
        last = st[last].fail;
      }
  st[last = cur].cnt += 1;
 void init(const char* s) {
  root = last = tot = 1;
  st[root] = node(0);
  for (char c; c = *s; ++s) extend(c - 'a');
 int q[N];
 void dp() +
  for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
  int head = 0, tail = 0;
  for (int i = 1; i <= tot; i++)</pre>
    if (st[i].indeg == 0) q[tail++] = i;
  while (head != tail) {
    int now = q[head++];
    if (int f = st[now].fail) {
      st[f].cnt += st[now].cnt;
      if (--st[f].indeg == 0) q[tail++] = f;
  }
 int run(const char* s) {
  int now = root;
  for (char c; c = *s; ++s)
    if (!st[now].ch[c -= 'a']) return 0;
    now = st[now].ch[c];
  return st[now].cnt;
 }
} SAM;
7.4 KMP
vector<int> kmp(const string &s) {
 vector<int> f(s.size(), 0);
 /* f[i] = length of the longest prefix
   (excluding s[0:i]) such that it coincides with the suffix of s[0:i] of the same length */
 /* i + 1 - f[i] is the length of the
   smallest recurring period of s[0:i] */
 int k = 0;
 for (int i = 1; i < (int)s.size(); ++i) {</pre>
  while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
  if (s[i] == s[k]) ++k;
  f[i] = k;
vector<int> search(const string &s, const string &t) {
 // return 0-indexed occurrence of t in s
 for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
  while(k > 0 && (k==(int)t.size() \mid \mid s[i]!=t[k]))
   k = f[k - 1];
  if (s[i] == t[k]) ++k;
```

if (k == (int)t.size()) r.push_back(i-t.size()+1);

7.5 Z value

| }

```
char s[MAXN];
int len,z[MAXN];
void Z_value() {
  int i,j,left,right;
  z[left=right=0]=len;
  for(i=1;i<len;i++) {
    j=max(min(z[i-left],right-i),0);
    for(;i+j<len&&s[i+j]==s[j];j++);
    if(i+(z[i]=j)>right)right=i+z[left=i];
  }
}
```

7.6 Manacher

```
int z[maxn];
int manacher(const string& s) {
    string t = ".";
    for(char c: s) t += c, t += '.';
    int l = 0, r = 0, ans = 0;
    for (int i = 1; i < t.length(); ++i) {
        z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
        while (i - z[i] >= 0 && i + z[i] < t.length()) {
        if(t[i - z[i]] == t[i + z[i]]) ++z[i];
        else break;
    }
    if (i + z[i] > r) r = i + z[i], l = i;
}
for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
    return ans;
}</pre>
```

7.7 Lexico Smallest Rotation

```
string mcp(string s){
  int n = s.length();
  s += s;
  int i=0, j=1;
  while (i<n && j<n){
    int k = 0;
    while (k < n && s[i+k] == s[j+k]) k++;
    if (s[i+k] <= s[j+k]) j += k+1;
    else i += k+1;
    if (i == j) j++;
}
int ans = i < n ? i : j;
  return s.substr(ans, n);
}</pre>
```

7.8 BWT

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a
 vector<int> v[ SIGMA ];
 void BWT(char* ori, char* res){
  // make ori -> ori + ori
  // then build suffix array
 }
 void iBWT(char* ori, char* res){
  for( int i = 0 ; i < SIGMA ; i ++ )</pre>
   v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )</pre>
   v[`ori[i] - BASE ].push_back( i´);
  vector<int> a;
  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
   for( auto j : v[ i ] ){
  a.push_back( j );
  ori[ ptr ++ ] = BASE + i;
  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
   ptr = a[ ptr ];
  res[ len ] = 0;
} bwt;
```

7.9 Palindromic Tree

```
struct palindromic_tree{
 struct node{
  int next[26],f,len;
  int cnt, num, st, ed;
  node( {\color{red} \textbf{int}} \ 1\text{=}\emptyset) : f(\emptyset), len(1), cnt(\emptyset), num(\emptyset) \ \{
   memset(next, 0, sizeof(next)); }
 vector<node> st:
 vector<char> s;
 int last.n:
 void init(){
  st.clear();s.clear();last=1; n=0;
  st.push_back(0);st.push_back(-1);
  st[0].f=1;s.push_back(-1); }
 int getFail(int x){
  while(s[n-st[x].len-1]!=s[n])x=st[x].f;
  return x;}
 void add(int c){
  s.push_back(c-='a'); ++n;
  int cur=getFail(last);
  if(!st[cur].next[c]){
   int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
   st[now].num=st[st[now].f].num+1;
  last=st[cur].next[c];
  ++st[last].cnt;}
 void dpcnt() {
  for (int i=st.size()-1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
 int size(){ return st.size()-2;}
} pt;
int main() {
 string s; cin >> s; pt.init();
 for (int i=0; i<SZ(s); i++) {</pre>
  int prvsz = pt.size(); pt.add(s[i]);
  if (prvsz != pt.size()) {
   int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [1,r]: s.substr(1, r-1+1)
 return 0;
```

8 Misc

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|{\rm det}(\tilde{L}_{rr})|.$

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\dots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \le k \le n$.

8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W\subseteq X, |W|\leq |N_G(W)|$

8.1.7 Euler's planar graph formula

```
V - E + F = C + 1, E \le 3V - 6(?)
```

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Lucas's theorem

```
 \binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p} \text{, where } m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,  and n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0.
```

8.1.10 Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2)$, find maximum $S\in I_1\cap I_2$. For each iteration, build the directed graph and find a shortest path from s to t.

- $s \to x : S \sqcup \{x\} \in I_1$ • $x \to t : S \sqcup \{x\} \in I_2$
- $y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1$ (y is in the unique circuit of $S \sqcup \{x\}$)
- $x \to y : S \setminus \{y\} \sqcup \{x\} \in I_2$ (y is in the unique circuit of $S \sqcup \{x\}$)

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|.$ In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

8.2 DP-opt Condition

8.2.1 totally monotone (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

8.2.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

8.3 Convex 1D/1D DP

```
struct segment {
int i, 1, r;
segment() {}
segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
dp[0] = 0;
deque<segment> dq; dq.push_back(segment(0, 1, n));
for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i);</pre>
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().l = i + 1
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
  f(i, dq.back().1) < f(dq.back().i, dq.back().1))
    dq.pop_back();
  if (dq.size())
   int d = 1 << 20, c = dq.back().1;</pre>
   while (d \gg 1) if (c + d \ll d, back().r)
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
  if (seg.1 <= n) dq.push_back(seg);</pre>
```

8.4 ConvexHull Optimization

```
struct Line {
  mutable int64_t a, b, p;
  bool operator<(const Line &rhs) const { return a < rhs
     .a; }
  bool operator<(int64_t x) const { return p < x; }
};
struct DynamicHull : multiset<Line, less<>> {
```

```
static const int64_t kInf = 1e18;
 bool Isect(iterator x, iterator y)
  auto Div = [](int64_t a, int64_t b) {
    return a / b - ((a ^ b) < 0 && a % b); }
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(int64_t a, int64_t b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
  if (x != begin() \&\& Isect(--x, y)) Isect(x, y = erase)
    (y))
  while ((y = x) != begin() \&\& (--x)->p >= y->p) Isect(
    x, erase(y));
 int64_t Query(int64_t x) {
  auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
};
```

8.5 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

8.6 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
 dfn[u]=low[u]=++dfs_idx;
 for(int i=0;i<(int)init_g[u].size();i++){</pre>
  int v=init_g[u][i];
  if(v==par[u]) continue;
  if(!dfn[v]){
   par[v]=u;
   tarjan(v);
   low[u]=min(low[u],low[v]);
   if(dfn[u]<low[v]){</pre>
    g[u].push_back(v)
    g[v].push_back(u);
  }else{
   low[u]=min(low[u],dfn[v]);
   if(dfn[v]<dfn[u]){
    int temp_v=u;
    bcc_id++
    while(temp_v!=v){
     g[bcc_id+n].push_back(temp_v);
     g[temp_v].push_back(bcc_id+n);
     temp_v=par[temp_v];
    g[bcc_id+n].push_back(v);
    g[v].push_back(bcc_id+n);
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
void dfs(int u,int fa){
 if(u<=n){
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
   dfs(v,u);
   memset(tp,0x8f,sizeof tp);
   if(v<=n){
```

tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);

tp[1]=max(

```
dp[u][0]+dp[v][0]+1
                                                                    ret.push_back( 2 );
     dp[u][1]+max(dp[v][0],dp[v][1])
   }else{
    tp[0]=dp[u][0]+dp[v][0];
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
   dp[u][0]=tp[0],dp[u][1]=tp[1];
 }else{
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
                                                                 }
   if(v==fa) continue;
   dfs(v,u);
  min_dp[0][0]=0;
min_dp[1][1]=1;
min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f3;
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
   memset(tmp,0x8f,sizeof tmp);
   tmp[0][0]=max(
    min_dp[0][0]+max(dp[v][0],dp[v][1]),
    min_dp[0][1]+dp[v][0]
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
   tmp[1][0]=max(
    min_dp[1][0]+max(dp[v][0],dp[v][1]),
    min_dp[1][1]+dp[v][0]
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
   memcpy(min_dp,tmp,sizeof tmp);
  dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
  dp[u][0]=min_dp[0][0];
int main(){
 int m,a,b;
 scanf("%d%d",&n,&m);
                                                                 }
 for(int i=0;i<m;i++){
  scanf("%d%d",&a,&b);</pre>
  init_g[a].push_back(b);
  init_g[b].push_back(a);
 par[1]=-1;
 tarjan(1);
 dfs(1,-1);
 printf("%d\n", max(dp[1][0], dp[1][1]));
 return 0;
8.7 Tree Knapsack
int dp[N][K]; PII obj[N];
vector<int> G[N];
void dfs(int u, int mx){
 for(int s: G[u]) {
  if(mx < obj[s].first) continue;</pre>
  for(int i=0;i<=mx-obj[s].FF;i++)</pre>
  dp[s][i] = dp[u][i];
dfs(s, mx - obj[s].first);
  for(int i=obj[s].FF;i<=mx;i++)</pre>
   dp[u][i] = max(dp[u][i],
    dp[s][i - obj[s].FF] + obj[s].SS);
}
8.8 N Queens Problem
vector< int > solve( int n ) {
 // no solution when n=2, 3
 vector< int > ret;
 if ( n % 6 == 2 ) {
 for ( int i = 2 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
  ret.push_back( 3 ); ret.push_back( 1 );
  for ( int i = 7 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 5 );
 } else if ( n % 6 == 3 ) {
for ( int i = 4 ; i <= n ; i += 2 )</pre>
   ret.push_back( i );
```

```
for ( int i = 5 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 1 ); ret.push_back( 3 );
 } else {
  for ( int i = 2 ; i <= n ; i += 2 )
   ret.push_back( i );
  for ( int i = 1 ; i <= n ; i += 2 )
   ret.push_back( i );
 return ret;
8.9 Aliens Optimization
long long Alien() {
 long long c = kInf;
 for (int d = 60; d >= 0; --d) {
  // cost can be negative, depending on the problem.
  if (c - (1LL << d) < 0) continue;</pre>
  long long ck = c - (1LL \ll d)
  pair<long long, int> r = check(ck);
  if (r.second == k) return r.first - ck * k;
  if (r.second < k) c = ck;
pair<long long, int> r = check(c);
 return r.first - c * k;
8.10 Hilbert Curve
long long hilbert(int n, int x, int y) {
 long long res = 0;
 for (int s = n / 2; s; s >>= 1) {
 int rx = (x & s) > 0, ry = (y & s) > 0;
res += s * 111 * s * ((3 * rx) ^ ry);
  if (ry == 0) {
   if (rx == 1) x = s - 1 - x, y = s - 1 - y;
   swap(x, y);
 return res;
8.11 Binary Search On Fraction
struct Q {
11 p, q;
 Q go(Q b, 11 d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(11 N) {
 Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
 assert(pred(hi));
 bool dir = 1, L = 1, H = 1;
 for (; L || H; dir = !dir) {
 11 len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)
   if (Q mid = hi.go(lo, len + step);
     mid.p > N \mid\mid mid.q > N \mid\mid dir ^ pred(mid))
    t++;
   else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
 return dir ? hi : lo;
```