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1.	1 '	vimr	С	

```
se is nu bs=2 ru mouse=a encoding=utf-8 ls=2
se cin cino+=j1 et sw=4 sts=4 tgc sc hls
syn on
colorscheme desert
filetype indent on
inoremap {<CR> {<CR>}<ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -std=c++17 -
    DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
    Wconversion -fsanitize=address,undefined -g && echo
     success<CR>
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 &&
     echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
```

1.2 Debug Macro

```
#ifdef KISEKI
#define safe cerr<<__PRETTY_FUNCTION__\
<<" line "<<__LINE__<<" safe\n"</pre>
#define debug(a...) qwerty(#a, a)
#define orange(a...) dvorak(#a, a)
using std::cerr;
template <typename ...T>
void qwerty(const char *s, T ...a) {
  cerr << "\e[1;32m(" << s << ") = (";
  int cnt = sizeof...(T);</pre>
   (..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
}
template <typename Iter>
void dvorak(const char *s, Iter L, Iter R) {
  cerr << "\e[1;32m[ " << s << " ] = [ ";</pre>
  for (int f = 0; L != R; ++L)
  cerr << (f++ ? ", " : "") << *L;
cerr << " ]\e[0m\n";</pre>
#else
#define safe ((void)0)
#define debug(...) ((void)0)
#define orange(...) ((void)0)
#endif
```

1.3 Increase Stack

```
const int size = 256 << 20;</pre>
                                    register long rsp asm("rsp");
                                    char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
                                    __asm__("movq %0, %%rsp\n"::"r"(p));
                                    // main
```

1.4 Pragma Optimization

```
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
1.5 IO Optimization
static inline int gc() {
  constexpr int B = 1 << 20;
  static char buf[B], *p, *q;
  if(p == a \&\&
    (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
   return EOF;
  return *p++;
template < typename T >
static inline bool gn( T &x ) {
 int c = gc(); T sgn = 1; x = 0;
while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
if(c == '-') sgn = -1, c = gc();
 if(c == EOF) return false;
 while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
 return x *= sgn, true;
```

#pragma GCC optimize("Ofast,no-stack-protector")

#pragma GCC optimize("no-math-errno,unroll-loops")

2 Data Structure

2.1 Dark Magic

2.2 Link-Cut Tree

p->ch[dir]=c;

```
struct Node{
Node *par, *ch[2];
int xor_sum, v;
bool is_rev;
Node(int _v){
 v=xor_sum=_v;is_rev=false;
 par=ch[0]=ch[1]=nullptr;
inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
inline void down(){
 if(is_rev){
  if(ch[0]!=nullptr) ch[0]->set_rev();
   if(ch[1]!=nullptr) ch[1]->set_rev();
   is_rev=false;
 }
inline void up(){
 xor_sum=v;
  if(ch[0]!=nullptr){
  xor_sum^=ch[0]->xor_sum;
  ch[0]->par=this;
 if(ch[1]!=nullptr){
  xor_sum^=ch[1]->xor_sum;
  ch[1]->par=this;
inline bool is_root(){
 return par==nullptr ||\
   (par->ch[0]!=this && par->ch[1]!=this);
bool is_rch(){return !is_root() && par->ch[1]==this;}
} *node[maxn], *stk[maxn];
int top;
void to_child(Node* p,Node* c,bool dir){
```

```
p->up();
inline void rotate(Node* node){
 Node* par=node->par;
 Node* par_par=par->par;
 bool dir=node->is_rch()
 bool par_dir=par->is_rch()
 to_child(par, node->ch[!dir], dir);
 to_child(node,par,!dir);
 if(par_par!=nullptr && par_par->ch[par_dir]==par)
  to_child(par_par,node,par_dir);
 else node->par=par_par;
inline void splay(Node* node){
 Node* tmp=node;
 stk[top++]=node;
 while(!tmp->is_root()){
  tmp=tmp->par;
  stk[top++]=tmp;
 while(top) stk[--top]->down();
 for(Node *fa=node->par;
  !node->is_root();
  rotate(node), fa=node->par)
  if(!fa->is_root())
   rotate(fa->is_rch()==node->is_rch()?fa:node);
inline void access(Node* node){
 Node* last=nullptr;
 while(node!=nullptr){
  splay(node);
  to_child(node, last, true);
  last=node;
  node=node->par;
inline void change_root(Node* node){
 access(node);splay(node);node->set_rev();
inline void link(Node* x, Node* y){
 change_root(x);splay(x);x->par=y;
inline void split(Node* x,Node* y){
 change_root(x);access(y);splay(x);
 to_child(x,nullptr,true);y->par=nullptr;
inline void change_val(Node* node,int v){
access(node);splay(node);node->v=v;node->up();
inline int query(Node* x,Node* y){
 change_root(x);access(y);splay(y);
 return y->xor_sum;
inline Node* find_root(Node* node){
 access(node);splay(node);
 Node* last=nullptr:
 while(node!=nullptr){
  node->down();last=node;node=node->ch[0];
 return last;
set<pii> dic;
inline void add_edge(int u,int v){
 if(u>v) swap(u,v)
 if(find_root(node[u])==find_root(node[v])) return;
 dic.insert(pii(u,v))
link(node[u],node[v]);
inline void del_edge(int u,int v){
 if(u>v) swap(u,v);
 if(dic.find(pii(u,v))==dic.end()) return;
 dic.erase(pii(u,v))
 split(node[u],node[v]);
2.3 LiChao Segment Tree
struct Line{
 int m, k, id;
 Line() : id( -1 ) {}
Line('int a, int'b,'int c')
: m(a), k(b), id(c) {}
```

int at(int x) { return m * x + k; }

```
private:
class LiChao {
                                                              vector< vector< T > > tbl;
private:
                                                              vector< int > lg;
                                                              T cv(Ta, Tb) {
  int n; vector< Line > nodes;
  inline int lc( int x ) { return 2 * x + 1; }
                                                               return Cmp_()( a, b ) ? a : b;
  inline int rc( int x ) { return 2 * x + 2; }
  void insert( int 1, int r, int id, Line ln ) {
                                                             public:
   int m = (1 + r) >> 1;
                                                              void init( T arr[], int n ) {
   if ( nodes[ id ].id == -1 ) {
                                                               // 0-base
   nodes[ id ] = ln;
                                                               lg.resize(n+1);
                                                               lg[0] = -1;
    return:
                                                               for( int i=1 ; i<=n ; ++i ) lg[i] = lg[i>>1] + 1;
   bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
                                                               tbl.resize(lg[n] + 1);
   if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
                                                               tbl[ 0 ].resize( n );
                                                               copy( arr, arr + n, tbl[ 0 ].begin() );
   atLeft ^= 1; swap( nodes[ id ], ln );
                                                               for ( int i = 1 ; i <= lg[ n ] ; ++ i ) {
  int len = 1 << ( i - 1 ), sz = 1 << i;</pre>
   if ( r - l == 1 ) return;
                                                                tbl[ i ].resize( n - sz + 1 );
   if ( atLeft ) insert( l, m, lc( id ), ln );
                                                                for (int_j = 0; j \le n - sz; ++ j
   else insert( m, r, rc( id ), ln );
                                                                 tbl[i][j] = cv(tbl[i-1][j], tbl[i-1][j+len]);
  int query( int 1, int r, int id, int x ) {
   int ret = 0;
   if ( nodes[ id ].id != -1 )
                                                              T query( int 1, int r ) {
                                                               // 0-base [1, r)
   ret = nodes[ id ].at( x );
   int m = (1 + r) >> 1;
                                                               int wh = lg[ r - l ], len = 1 << wh;</pre>
                                                               return cv( tbl[ wh ][ 1 ], tbl[ wh ][ r - len ] );
   if ( r - l == 1 ) return ret;
   else if (x < m )
                                                            };
    return max( ret, query( 1, m, lc( id ), x ) );
   else
                                                             2.6
                                                                  Linear Basis
    return max( ret, query( m, r, rc( id ), x ) );
                                                             template <int BITS>
public:
                                                             struct LinearBasis {
 void build( int n_ ) {
                                                              array<uint64_t, BITS> basis;
                                                              Basis() { basis.fill(0); }
  n = n_; nodes.clear();
  nodes.resize( n << 2, Line() );</pre>
                                                              void add(uint64_t x)
                                                               for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
 void insert( Line ln ) { insert( 0, n, 0, ln ); }
                                                                if (basis[i] == 0) {
  int query( int x ) { return query( 0, n, 0, x ); }
                                                                 basis[i] = x;
} lichao;
                                                                 return;
2.4 Treap
                                                                x ^= basis[i];
namespace Treap{
                                                               }
#define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                              bool ok(uint64_t x) {
struct node{
                                                               for (int i = 0; i < BITS; ++i)</pre>
 int size;
  uint32_t pri;
                                                                if ((x >> i) & 1) x ^= basis[i];
 node *lc, *rc;
node() : size(0), pri(rand()), lc(0), rc(0) {}
                                                               return x == 0;
                                                            }:
  void pull() {
  size = 1;
                                                                   Binary Search On Segment Tree
  if ( lc ) size += lc->size;
   if ( rc ) size += rc->size;
                                                            // find_first = x -> minimal x s.t. check( [a, x) )
                                                             // find_last = x \rightarrow \max x s.t. check([x, b))
  }
                                                             template <typename C>
node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
                                                             int find_first(int 1, const C &check) {
                                                              if (1 >= n)
 if ( L->pri > R->pri ) {
                                                               return n;
  L->rc = merge( L->rc, R ); L->pull();
                                                              1 += sz;
   return L;
                                                              for (int i = height; i > 0; i--)
 } else {
                                                               propagate(1 >> i);
  R->lc = merge( L, R->lc ); R->pull();
                                                              Monoid sum = identity;
                                                              do {
   return R;
                                                               while ((1 \& 1) == 0)
}
                                                                1 >>= 1
void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                               if (check(f(sum, data[1]))) {
 if ( not rt ) L = R = nullptr;
                                                                while (1 < sz) {</pre>
  else if( sz( rt->lc ) + 1 <= k ) {
                                                                 propagate(1);
  split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                 auto nxt = f(sum, data[1]);
  L->pull();
                                                                 if (not check(nxt)) {
 } else {
                                                                  sum = nxt;
  R = rt:
                                                                  1++;
   split_by_size( rt->lc, k, L, R->lc );
   R->pull();
 }
                                                                return 1 + 1 - sz;
#undef sz
                                                               sum = f(sum, data[1++]);
                                                              } while ((1 & -1) != 1);
                                                              return n;
2.5 Sparse Table
template < typename T, typename Cmp_ = less< T > >
class SparseTable {
                                                            template <typename C>
```

if (not ins[t]) {

st.push_back(t);

```
int find_last(int r, const C &check) {
                                                                  ins[t] = true;
if (r <= 0)
                                                                 if (dfn[v]) {
 return -1;
                                                                  low[u] = min(low[u], dfn[v]);
 r += sz;
for (int i = height; i > 0; i--)
                                                                  continue
 propagate((r - 1) >> i);
                                                                 } ++ch; dfs(v, u)
Monoid sum = identity;
                                                                 low[u] = min(low[u], low[v]);
                                                                 if (low[v] >= dfn[u]) {
do {
                                                                  ap[u] = true;
 while (r > 1 \text{ and } (r \& 1))
                                                                  while (true) {
                                                                   int eid = st.back(); st.pop_back();
  r >>= 1
  if (check(f(data[r], sum))) {
                                                                   bcc[eid] = ecnt;
                                                                   if (eid == t) break;
  while (r < sz) {</pre>
   propagate(r);
    r = (r << 1) + 1;
                                                                  ecnt++;
   auto nxt = f(data[r], sum);
    if (not check(nxt)) {
    sum = nxt;
                                                                if (ch == 1 and u == f) ap[u] = false;
                                                               }
     r--;
   }
                                                              public:
                                                               void init(int n_) {
   return r - sz;
                                                                G.clear(); G.resize(n = n_);
                                                                ecnt = 0; ap.assign(n, false);
 sum = f(data[r], sum);
                                                                low.assign(n, 0); dfn.assign(n, 0);
} while ((r & -r) != r);
return -1;
                                                               void add_edge(int u, int v) {
                                                               G[u].emplace_back(v, ecnt);
                                                                G[v].emplace_back(u, ecnt++);
3
    Graph
                                                               void solve() {
                                                                ins.assign(ecnt, false);
   BCC Edge
                                                                bcc.resize(ecnt); ecnt = 0;
class BCC_Bridge {
                                                                for (int i = 0; i < n; ++i)</pre>
private:
                                                                 if (not dfn[i]) dfs(i, i);
 int n. ecnt:
 vector<vector<pair<int,int>>> G;
                                                               int get_id(int x) { return bcc[x]; }
  vector<int> dfn, low;
                                                               int count() { return ecnt; }
 vector<bool> bridge;
                                                               bool is_ap(int x) { return ap[x]; }
  void dfs(int u, int f)
                                                            } bcc_ap;
  dfn[u] = low[u] = dfn[f] + 1;
  for (auto [v, t]: G[u]) {
                                                             3.3 2-SAT (SCC)
    if (v == f) continue;
                                                            class TwoSat{
    if (dfn[v]) {
                                                              private:
    low[u] = min(low[u], dfn[v]);
                                                               int n:
     continue;
                                                               vector<vector<int>> rG,G,sccs;
                                                               vector<int> ord,idx;
   dfs(v, u);
                                                               vector<bool> vis,result;
   low[u] = min(low[u], low[v]);
                                                               void dfs(int u){
    if (low[v] > dfn[u]) bridge[t] = true;
                                                                vis[u]=true
                                                                for(int v:G[u])
                                                                 if(!vis[v]) dfs(v);
public:
                                                                ord.push_back(u);
  void init(int n_) {
  G.clear(); G.resize(n = n_);
                                                               void rdfs(int u){
  low.assign(n, ecnt = 0);
                                                                vis[u]=false;idx[u]=sccs.size()-1;
  dfn.assign(n, 0);
                                                                sccs.back().push_back(u);
                                                                for(int v:rG[u])
  void add_edge(int u, int v) {
                                                                 if(vis[v])rdfs(v);
  G[u].emplace_back(v, ecnt);
G[v].emplace_back(u, ecnt++);
                                                              public:
                                                               void init(int n_){
 void solve() {
                                                                n=n_;G.clear();G.resize(n);
  bridge.assign(ecnt, false);
                                                                rG.clear();rG.resize(n);
  for (int i = 0; i < n; ++i)</pre>
                                                                sccs.clear();ord.clear();
   if (not dfn[i]) dfs(i, i);
                                                                idx.resize(n);result.resize(n);
 bool is_bridge(int x) { return bridge[x]; }
                                                               void add_edge(int u,int v){
} bcc_bridge;
                                                                G[u].push_back(v);rG[v].push_back(u);
3.2 BCC Vertex
                                                               void orr(int x,int y){
class BCC_AP {
                                                                if ((x^y)==1) return
                                                                add_edge(x^1,y); add_edge(y^1,x);
private:
 int n, ecnt;
                                                               bool solve(){
 vector<vector<pair<int,int>>> G;
 vector<int> bcc, dfn, low, st;
                                                                vis.clear();vis.resize(n);
 vector<bool> ap, ins;
                                                                for(int i=0;i<n;++i)</pre>
 void dfs(int u, int f) {
  dfn[u] = low[u] = dfn[f] + 1;
                                                                 if(not vis[i])dfs(i);
                                                                reverse(ord.begin(),ord.end());
  int ch = 0;
                                                                for (int u:ord){
   for (auto [v, t]: G[u]) if (v != f) {
                                                                 if(!vis[u])continue;
```

sccs.push_back(vector<int>());

rdfs(u);

time_ = 0; dfschain(1, 1);

```
for(int i=0;i<n;i+=2)</pre>
                                                                PII get_subtree(int u) { return {tl[ u ],tr[ u ] }; }
    if(idx[i]==idx[i+1])
                                                                vector< PII > get_path( int u , int v ){
                                                                 vector< PII > res;
     return false
                                                                 int g = lca( u, v );
   vector<bool> c(sccs.size());
   for(size_t i=0;i<sccs.size();++i){</pre>
                                                                 while ( chain[ u ] != chain[ g ] ) {
                                                                  int s = chain_st[ chain[ u ] ];
res.emplace_back( tl[ s ], tl[ u ] + 1 );
    for(size_t j=0;j<sccs[i].size();++j){</pre>
     result[sccs[i][j]]=c[i];
     c[idx[sccs[i][j]^1]]=!c[i];
                                                                  u = fa[ s ][ 0 ];
                                                                 res.emplace\_back( tl[ g ], tl[ u ] + 1 );
                                                                 while ( chain[ v ] != chain[ g ] ) {
   return true;
                                                                  int s = chain_st[ chain[ v ] ]
 bool get(int x){return result[x];}
                                                                  res.emplace_back( tl[ s ], tl[ v ] + 1 );
  inline int get_id(int x){return idx[x];}
                                                                  v = fa[ s ][ 0 ];
  inline int count(){return sccs.size();}
                                                                 res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
                                                                 return res:
3.4 Lowbit Decomposition
                                                                 /* res : list of intervals from u to v
class LowbitDecomp{
                                                                  * ( note only nodes work, not edge )
                                                                  * usage :
private:
                                                                  * vector< PII >& path = tree.get_path( u , v )
int time_, chain_, LOG_N;
                                                                  * for( auto [ 1, r ] : path ) {
vector< vector< int > > G, fa;
vector< int > tl, tr, chain, chain_st;
                                                                  * 0-base [ 1, r )
                                                                  * }
// chain_ : number of chain
                                                                  */
// tl, tr[ u ] : subtree interval in the seq. of u
                                                                }
 // chain_st[ u ] : head of the chain contains u
// chian[ u ] : chain id of the chain u is on
                                                              } tree;
void predfs( int u, int f ) {
                                                               3.5
                                                                     MaxClique
  chain[u] = 0;
  for ( int v : G[ u ] ) {
                                                               // contain a self loop u to u, than u won't in clique
  if ( v == f ) continue;
                                                               template < size_t MAXN >
   predfs( v, u );
                                                               class MaxClique{
   if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
                                                               private:
    chain[ u ] = chain[ v ];
                                                                using bits = bitset< MAXN >;
                                                                bits popped, G[ MAXN ], ans;
size_t deg[ MAXN ], deo[ MAXN ], n;
 if ( not chain[ u ] )
   chain[ u ] = chain_ ++;
                                                                void sort_by_degree() {
                                                                 popped.reset();
 void dfschain( int u, int f ) {
                                                                 for ( size_t i = 0 ; i < n ; ++ i )</pre>
 fa[ u ][ 0 ] = f;
for ( int i = 1 ; i < LOG_N ; ++ i )
                                                                   deg[ i ] = G[ i ].count();
                                                                 for ( size_t i = 0 ; i < n ; ++ i ) {
  fa[u][i] = fa[fa[u][i-1]][i-1];
                                                                   size_t mi = MAXN, id = 0;
                                                                   for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )
    mi = deg[ id = j ];</pre>
  tl[ u ] = time_++;
  if ( not chain_st[ chain[ u ] ] )
   chain_st[ chain[ u ] ] = u;
  for ( int v : G[ u ] )
                                                                   popped[ deo[ i ] = id ] = 1;
  if ( v != f and chain[ v ] == chain[ u ] )
                                                                   for( size_t u = G[ i ]._Find_first() ;
                                                                    u < n ; u = G[ i ]._Find_next( u ) )</pre>
    dfschain( v, u );
  for ( int v : G[ u ] )
                                                                     -- deg[ u ];
   if ( v != f and chain[ v ] != chain[ u ] )
    dfschain( v, u );
                                                                void BK( bits R, bits P, bits X ) {
  tr[ u ] = time_;
                                                                 if (R.count()+P.count() <= ans.count()) return;</pre>
bool anc( int u, int v ) {
  return tl[ u ] <= tl[ v ] and tr[ v ] <= tr[ u ];</pre>
                                                                 if ( not P.count() and not X.count() ) {
                                                                  if ( R.count() > ans.count() ) ans = R;
                                                                  return;
public:
                                                                 }
int lca( int u, int v ) {
                                                                 /* greedily chosse max degree as pivot
  if ( anc( u, v ) ) return u;
                                                                 bits cur = P | X; size_t pivot = 0, sz = 0;
 for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
if ( not anc( fa[ u ][ i ], v ) )
                                                                 for ( size_t u = cur._Find_first() ;
                                                                  u < n ; u = cur._Find_next( u )</pre>
   u = fa[ u ][ i ];
                                                                   if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
                                                                 cur = P & ( ~G[ pivot ] );
  return fa[ u ][ 0 ];
                                                                 */ // or simply choose first
void init( int n ) {
                                                                 bits cur = P & (~G[ ( P | X )._Find_first() ]);
 fa.assign( ++n, vector< int >( LOG_N ) );
                                                                 for ( size_t u = cur._Find_first()
                                                                  u < n ; u = cur._Find_next( u ) ) {
if ( R[ u ] ) continue;</pre>
  for ( LOG_N = 0 ; ( 1 << LOG_N ) < n ; ++ LOG_N );
 G.clear(); G.resize( n );
 tl.assign( n, 0 ); tr.assign( n, 0 );
                                                                  R[u] = 1;
 chain.assig( n, 0 ); chain_st.assign( n, 0 );
                                                                  BK( R, P & G[ u ], X & G[ u ] );
                                                                  R[u] = P[u] = 0, X[u] = 1;
 void add_edge( int u , int v ) {
  // 1-base
 G[ u ].push_back( v );
                                                               public:
 G[ v ].push_back( u );
                                                                void init( size_t n_ ) {
                                                                 n = n_{-};
void decompose(){
                                                                 for ( size_t i = 0 ; i < n ; ++ i )
 chain_ = 1;
                                                                  G[ i ].reset();
 predfs( 1, 1 );
                                                                 ans.reset();
```

void add_edges(int u, bits S) { G[u] = S; }

```
void add_edge( int u, int v ) {
                                                                sort(r.begin(), r.end(),
                                                                 [&](int i, int j) { return d[i] > d[j]; });
  G[u][v] = G[v][u] = 1;
                                                                csort(r, c);
                                                                dfs(r, c, 1, mask);
return ans; // sol[0 ~ ans-1]
 int solve() {
  sort_by_degree(); // or simply iota( deo... )
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                               }
   deg[ i ] = G[ i ].count();
                                                              } graph;
  bits pob, nob = 0; pob.set();
                                                              3.7 Virtural Tree
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
  for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t v = deo[ i ];</pre>
                                                              inline bool cmp(const int &i, const int &j) {
                                                               return dfn[i] < dfn[j];</pre>
   bits tmp; tmp[ v ] = 1;
   BK( tmp, pob & G[ v ], nob & G[ v ] );
pob[ v ] = 0, nob[ v ] = 1;
                                                              void build(int vectrices[], int k) {
                                                               static int stk[MAX_N];
                                                               sort(vectrices, vectrices + k, cmp);
  return static_cast< int >( ans.count() );
                                                               stk[sz++] = 0;
                                                               for (int i = 0; i < k; ++i) {
  int u = vectrices[i], lca = LCA(u, stk[sz - 1]);</pre>
};
                                                                if (lca == stk[sz - 1]) stk[sz++] = u;
3.6 MaxCliqueDyn
                                                                else {
constexpr int kN = 150;
                                                                 while (sz \ge 2 \&\& dep[stk[sz - 2]] \ge dep[lca]) {
struct MaxClique { // Maximum Clique
                                                                  addEdge(stk[sz - 2], stk[sz - 1]);
sz--:
                                                                 if (stk[sz - 1] != lca) {
 void init(int _n) {
 n = n, ans q = 0;
                                                                  addEdge(lca, stk[--sz]);
  for (int i = 0; i < n; i++) a[i].reset();</pre>
                                                                  stk[sz++] = lca, vectrices[cnt++] = lca;
 void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
void csort(vector<int> &r, vector<int> &c) {
                                                                 stk[sz++] = u;
 int mx = 1, km = max(ans - q + 1, 1), t = 0,
                                                               for (int i = 0; i < sz - 1; ++i)
    m = int(r.size())
  cs[1].reset(); cs[2].reset()
                                                                addEdge(stk[i], stk[i + 1]);
  for (int i = 0; i < m; i++) {
   int p = r[i], k = 1
                                                              3.8 Centroid Decomposition
   while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
                                                              struct Centroid {
   cs[k][p] = 1;
                                                               vector<vector<int64_t>> Dist;
                                                               vector<int> Parent, Depth;
   if (k < km) r[t++] = p;
                                                               vector<int64_t> Sub, Sub2;
  c.resize(m);
                                                               vector<int> Sz, Sz2;
  if(t) c[t-1] = 0;
                                                               Centroid(vector<vector<pair<int, int>>> g) {
  for (int k = km; k <= mx; k++) {</pre>
                                                                int N = g.size()
  for (int p = int(cs[k]._Find_first());
                                                                vector<bool> Vis(N);
      p < kN; p = int(cs[k]._Find_next(p))) {
                                                                vector<int> sz(N), mx(N);
    r[t] = p; c[t++] = k;
                                                                vector<int> Path;
                                                                Dist.resize(N)
  }
                                                                Parent.resize(N);
                                                                Depth.resize(N)
 void dfs(vector<int> &r, vector<int> &c, int 1,
                                                                auto DfsSz = [\&](auto dfs, int x) -> void {
  bitset<kN> mask) {
                                                                 Vis[x] = true; sz[x] = 1; mx[x] = 0;
                                                                 for (auto [u, w] : g[x]) {
  while (!r.empty()) {
                                                                  if (Vis[u]) continue;
   int p = r.back(); r.pop_back();
                                                                  dfs(dfs, u)
   mask[p] = 0;
   if (q + c.back() <= ans) return;</pre>
                                                                  sz[x] += sz[u];
                                                                  mx[x] = max(mx[x], sz[u]);
   cur[q++] = p;
   vector<int> nr, nc;
   bitset<kN> nmask = mask & a[p];
                                                                 Path.push_back(x);
   for (int i : r)
                                                                };
    if (a[p][i]) nr.push_back(i);
                                                                auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
                                                                 -> void {
   if (!nr.empty()) {
    if (1 < 4) {
                                                                 Dist[x].push_back(D); Vis[x] = true;
     for (int i : nr)
                                                                 for (auto [u, w] : g[x]) {
      d[i] = int((a[i] & nmask).count());
                                                                  if (Vis[u]) continue;
     sort(nr.begin(), nr.end(),
                                                                  dfs(dfs, u, D + w);
      [&](int x, int y)
                                                                 }
       return d[x] > d[y];
                                                                };
                                                                auto Dfs = [&]
      });
                                                                 (auto dfs, int x, int D = 0, int p = -1)->void {
   csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
                                                                 Path.clear(); DfsSz(DfsSz, x);
                                                                 int M = Path.size();
                                                                 int C = -1;
    ans = q; copy(cur, cur + q, sol);
                                                                 for (int u : Path) {
   c.pop_back(); q--;
                                                                  if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
  }
                                                                  Vis[u] = false;
 int solve(bitset<kN> mask) { // vertex mask
                                                                 DfsDist(DfsDist, C);
                                                                 for (int u : Path) Vis[u] = false;
  vector<int> r, c;
  for (int i = 0; i < n; i++)
                                                                 Parent[C] = p; Vis[C] = true;
  if (mask[i]) r.push_back(i);
for (int i = 0; i < n; i++)</pre>
                                                                 Depth[C] = D;
                                                                 for (auto [u, w] : g[C]) {
                                                                  if (Vis[u]) continue;
   d[i] = int((a[i] & mask).count());
```

```
dfs(dfs, u, D + 1, C);
                                                                  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
   }
                                                                  for (int i=n; !vst[st]; st=prv[i--][st]) {
  Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
                                                                   vst[st]++
                                                                   edgeID.PB(prve[i][st]);
  Sz.resize(N); Sz2.resize(N);
                                                                   rho.PB(st);
 void Mark(int v) {
  int x = v, z = -1;
                                                                  while (vst[st] != 2) {
  for (int i = Depth[v]; i >= 0; --i) {
                                                                   int v = rho.back(); rho.pop_back();
                                                                   cycle.PB(v);
   Sub[x] += Dist[v][i]; Sz[x]++;
   if (z != -1) {
                                                                   vst[v]++;
    Sub2[z] += Dist[v][i];
    Sz2[z]++;
                                                                  reverse(ALL(edgeID));
                                                                  edgeID.resize(SZ(cycle));
   z = x; x = Parent[x];
                                                                  return mmc;
  }
                                                               } mmc;
 int64_t Query(int v) {
                                                                3.11 Mo's Algorithm on Tree
 int64_t res = 0;
  int x = v, z = -1
                                                                int q; vector< int > G[N];
 for (int i = Depth[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
                                                                struct Que{
                                                                int u, v, id;
} que[ N ];
   if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
                                                                int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
   z = x; x = Parent[x];
                                                                void_dfs( int u, int f ) {
                                                                 dfn[ u ] = dfn_++; int saved_rbp = stk_;
for ( int v : G[ u ] ) {
  return res;
                                                                  if ( v == f ) continue;
};
                                                                  dfs( v, u );
3.9
     Tree Hashing
                                                                  if ( stk_ - saved_rbp < SQRT_N ) continue;</pre>
                                                                  for ( ++ block_ ; stk_ != saved_rbp ; )
  block_id[ stk[ -- stk_ ] ] = block_;
uint64_t hsah(int u, int f) {
 uint64_t r = 127;
 for (int v : G[ u ]) if (v != f) {
 uint64_t hh = hsah(v, u);
                                                                stk[ stk_ ++ ] = u;
  r=(r+(hh*hh)%1010101333)%1011820613;
                                                                bool inPath[ N ];
                                                                void Diff( int u ) {
return r;
                                                                if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
                                                                 else { /*add this edge*/ }
3.10 Minimum Mean Cycle
/* minimum mean cycle O(VE) */
                                                                void traverse( int& origin_u, int u ) {
                                                                for ( int g = lca( origin_u, u )
struct MMC{
                                                                  origin_u != g ; origin_u = parent_of[ origin_u ] )
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
                                                                   Diff( origin_u );
#define V 1021
                                                                 for (int v = u; v != origin_u; v = parent_of[v])
#define inf 1e9
                                                                  Diff( v );
 struct Edge { int v,u; double c; };
                                                                 origin_u = u;
 int n, m, prv[V][V], prve[V][V], vst[V];
                                                                }
 Edge e[E];
                                                                void solve() {
 vector<int> edgeID, cycle, rho;
                                                                 dfs( 1, 1 );
 double d[V][V];
                                                                 while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
 void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
                                                                 sort( que, que + q, [](const Que& x, const Que& y) {
                                                                  return tie( block_id[ x.u ], dfn[ x.v ] )
 void add_edge( int vi , int ui , double ci )
                                                                      < tie( block_id[ y.u ], dfn[ y.v ] );
 { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
                                                                 } );
                                                                 int U = 1, V = 1;
                                                                 for ( int i = 0 ; i < q ; ++ i ) {
  for(int i=0; i<n; i++) d[0][i]=0;
  for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);</pre>
                                                                  pass( U, que[ i ].u );
pass( V, que[ i ].v );
   for(int j=0; j<m; j++) {</pre>
                                                                  // we could get our answer of que[ i ].id
    int v = e[j].v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
                                                               }
                                                                /*
     d[i+1][u] = d[i][v]+e[j].c;
     prv[i+1][u] = v;
                                                                Method 2:
     prve[i+1][u] = j;
                                                               dfs u:
                                                                push u
                                                                 iterate subtree
                                                                Let P = LCA(u, v), and St(u) \le St(v)
                                                               if (P == u) query[St(u), St(v)]
 double solve(){
  // returns inf if no cycle, mmc otherwise
                                                                else query[Ed(u), St(v)], query[St(P), St(P)]
  double mmc=inf;
  int st = -1
                                                                3.12 Minimum Steiner Tree
  bellman_ford();
  for(int i=0; i<n; i++) {</pre>
                                                               // Minimum Steiner Tree
   double avg=-inf;
                                                               // 0(V 3^T + V^2 2^T)
                                                               struct SteinerTree{
   for(int k=0; k<n; k++) {</pre>
    if(d[n][i]<inf-eps)</pre>
                                                               #define V 33
     avg=max(avg,(d[n][i]-d[k][i])/(n-k));
                                                                #define T 8
                                                               #define INF 1023456789
    else avg=max(avg,inf);
                                                                 int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
                                                                void init( int _n ){
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
```

if (in[i] == -inf)
return -inf;

// find cycle

```
int tot = 0;
  for( int i = 0 ; i < n ; i ++ ){</pre>
                                                                  vector<int> id(n, -1), vis(n, -1);
   for( int j = 0 ; j < n ; j ++ )</pre>
                                                                   for (int i = 0; i < n; i++) {
   dst[ i ][ j ] = INF;
dst[ i ][ i ] = 0;
                                                                    ans += in[i];
                                                                    for (int x = i; x != -1 && id[x] == -1; x = prv[x])
  }
                                                                     if (vis[x] == i) {
                                                                     for (int y = prv[x]; y != x; y = prv[y])
 void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
                                                                       id[y] = tot;
  dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
                                                                      id[x] = tot++;
                                                                      break:
 void shortest_path(){
  for( int k = 0 ; k < n ; k ++ )</pre>
                                                                     vis[x] = i;
   for( int i = 0 ; i < n ; i ++ )</pre>
                                                                   }
    for( int j = 0 ; j < n ; j ++ )</pre>
     dst[ i ][ j ] = min( dst[ i ][ j ],
    dst[ i ][ k ] + dst[ k ][ j ] );
                                                                  if (!tot)
                                                                    return ans;
                                                                   for (int i = 0; i < n; i++)
 int solve( const vector<int>& ter ){
                                                                   if (id[i] == -1)
  int t = (int)ter.size();
                                                                     id[i] = tot++;
  for( int i = 0 ; i < ( 1 << t ) ; i ++ )
                                                                   // shrink
   for( int j = 0; j < n; j ++ )
dp[ i ][ j ] = INF;
                                                                  for (auto &e : E) {
  if (id[e.u] != id[e.v])
  for( int i = 0 ; i < n ; i ++ )
                                                                    e.w -= in[e.v];
   dp[0][i] = 0;
                                                                   e.u = id[e.u], e.v = id[e.v];
  for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){
  if( msk == ( msk & (-msk) ) ){</pre>
                                                                  n = tot:
    int who = __lg( msk );
                                                                  root = id[root];
    for( int i = 0 ; i < n ; i ++ )
dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
                                                                 assert(false);
    continue:
                                                               } DMST;
   for( int i = 0 ; i < n ; i ++ )</pre>
                                                               3.14
                                                                      Dominator Tree
    for( int submsk = ( msk - 1 ) & msk ; submsk ;
         submsk = ( submsk - 1 ) & msk )
                                                               namespace dominator {
      vector<int> g[maxn], r[maxn], rdom[maxn];
                                                               int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
               dp[ msk ^ submsk ][ i ] );
                                                               int dom[maxn], val[maxn], rp[maxn], tk;
   for( int i = 0 ; i < n ; i ++ ){</pre>
                                                               void init(int n) {
    tdst[ i ] = INF;
                                                                // vertices are numbered from 0 to n - 1
    fill(dfn, dfn + n, -1);fill(rev, rev + n, -1);
                                                                fill(fa, fa + n, -1); fill(val, val + n, -1);
                                                                fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
                                                                fill(dom, dom + n, -1); tk = \hat{0};
   for( int i = 0 ; i < n ; i ++ )</pre>
                                                                for (int i = 0; i < n; ++i) {
    dp[ msk ][ i ] = tdst[ i ];
                                                                 g[i].clear(); r[i].clear(); rdom[i].clear();
  int ans = INF:
  for( int i = 0 ; i < n ; i ++ )</pre>
                                                               void add_edge(int x, int y) { g[x].push_back(y); }
   ans = min( ans , dp[ (1 << t) - 1 ][ i ] );
                                                               void dfs(int x) {
                                                                rev[dfn[x] = tk] = x;
  return ans;
                                                                fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
} solver;
                                                                for (int u : g[x]) {
                                                                 if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
      Directed Minimum Spanning Tree
                                                                 r[dfn[u]].push_back(dfn[x]);
struct DirectedMST { // find maximum
 struct Edge {
  int u, v;
                                                               void merge(int x, int y) { fa[x] = y; }
                                                               int find(int x, int c = 0) {
  int w:
                                                                if (fa[x] == x) return c ? -1 : x;
  Edge(int u, int v, int w) : u(u), v(v), w(w) {}
                                                                int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
 vector<Edge> Edges;
 void clear() { Edges.clear(); }
                                                                if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
 void addEdge(int a, int b, int w) { Edges.emplace_back
                                                                fa[x] = p;
return c ? p : val[x];
    (a, b, w); }
 int solve(int root, int n) {
                                                               vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
  vector<Edge> E = Edges;
  int ans = 0;
                                                               // p[i] = -2 if i is unreachable from s
  while (true) {
   // find best in edge
                                                                dfs(s);
   vector<int> in(n, -inf), prv(n, -1);
                                                                for (int i = tk - 1; i >= 0; --i) {
  for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
                                                                 if (i) rdom[sdom[i]].push_back(i);
                                                                 for (int &u : rdom[i]) {
     in[e.v] = e.w;
                                                                  int p = find(u);
     prv[e.v] = e.u;
                                                                  if (sdom[p] == i) dom[u] = i;
else dom[u] = p;
   in[root] = 0;
   prv[root] = -1;
   for (int i = 0; i < n; i++)
                                                                 if (i) merge(i, rp[i]);
```

vector<int> p(n, -2); p[s] = -1;

for (int i = 1; i < tk; ++i)

```
if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
 for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
 return p;
3.15
      Edge Coloring
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
 for (int i = 0; i <= N; i++)
  for (int j = 0; j \leftarrow N; j++)
    C[i][j] = G[i][j] = 0;
void solve(vector<pair<int, int>> &E, int N) {
int X[kN] = {}, a;
auto update = [&](int u) {
  for (X[u] = 1; C[u][X[u]]; X[u]++);
 auto color = [&](int u, int v, int c) {
  int p = G[u][v];
  G[u][v] = G[v][u] = c;
  C[u][c] = v, C[v][c] = u;
  C[u][p] = C[v][p] = 0;
  if (p) X[u] = X[v] = p;
  else update(u), update(v);
  return p;
 } ;
 auto flip = [&](int u, int c1, int c2) {
  int p = C[u][c1];
  swap(C[u][c1], C[u][c2]);
 if (p) G[u][p] = G[p][u] = c2;
if (!C[u][c1]) X[u] = c1;
  if (!C[u][c2]) X[u] = c2;
  return p;
 };
 for (int i = 1; i <= N; i++) X[i] = 1;
 for (int t = 0; t < E.size(); t++) {</pre>
  auto [u, v] = E[t];
  int v0 = v, c = X[u], c0 = c, d;
  vector<pair<int, int>> L; int vst[kN] = {};
  while (!G[u][v0]) {
   L.emplace_back(v, d = X[v]);
   if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
   c = color(u, L[a].first, c);
else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
     color(u, L[a].first, L[a].second);
   else if (vst[d]) break
   else vst[d] = 1, v = C[u][d];
  if (!G[u][v0]) {
   for (; v; v = flip(v, c, d), swap(c, d));
   if (C[u][c0]) { a = int(L.size()) - 1;
    while (--a >= 0 && L[a].second != c);
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
   } else t--;
     Matching & Flow
4
    Kuhn Munkres
class KM {
private:
 static constexpr lld INF = 1LL << 60;</pre>
 vector<lld> h1,hr,slk;
```

```
private:
    static constexpr lld INF = 1LL << 60;
    vector<lld> hl,hr,slk;
    vector<int> fl,fr,pre,qu;
    vector<vector<lld> w;
    vector<bool> vl,vr;
    int n, ql, qr;
    bool check(int x) {
        if (vl[x] = true, fl[x] != -1)
            return vr[qu[qr++] = fl[x]] = true;
        while (x != -1) swap(x, fr[fl[x] = pre[x]]);
        return false;
    }
    void bfs(int s) {
        fill(slk.begin(), slk.end(), INF);
        fill(vl.begin(), vl.end(), false);
        fill(vr.begin(), vr.end(), false);
        ql = qr = 0;
```

```
while (true) {
   11d d;
   while (ql < qr) {</pre>
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
     if(!vl[x]&&slk[x]>=(d=hl[x]+hr[y]-w[x][y])){
      if (pre[x] = y, d) slk[x] = d;
       else if (!check(x)) return;
    }
   d = INF;
   for (int x = 0; x < n; ++x)
    if (!v1[x] \&\& d > s1k[x]) d = s1k[x];
   for (int x = 0; x < n; ++x) {
    if (v1[x]) h1[x] += d;
    else slk[x] -= d;
    if (vr[x]) hr[x] -= d;
   for (int x = 0; x < n; ++x)
    if (!v1[x] && !slk[x] && !check(x)) return;
public:
 void init( int n_ ) {
  n = n_; qu.resize(n);
  fl.clear(); fl.resize(n, -1);
fr.clear(); fr.resize(n, -1);
hr.clear(); hr.resize(n); hl.resize(n);
  w.clear(); w.resize(n, vector<lld>(n));
  slk.resize(n); pre.resize(n);
  vl.resize(n); vr.resize(n);
 void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 1ld solve() {
  for (int i = 0; i < n; ++i)</pre>
   hl[i] = *max_element(w[i].begin(), w[i].end());
  for (int i = 0; i < n; ++i) bfs(i);</pre>
  11d res = 0;
  for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
  return res;
 }
} km;
4.2 Bipartite Matching
class BipartiteMatching{
private:
 vector<int> X[N], Y[N];
 int fX[N], fY[N], n;
 bitset<N> walked;
 bool dfs(int x)
  for(auto i:X[x]){
   if(walked[i])continue;
   walked[i]=1;
   if(fY[i]==-1||dfs(fY[i])){
    fY[i]=x;fX[x]=i;
    return 1;
   }
  return 0;
public:
 void init(int _n){
  n=_n; walked.reset();
  for(int i=0;i<n;i++){</pre>
   X[i].clear();Y[i].clear();
   fX[i]=fY[i]=-1;
  }
 void add_edge(int x, int y){
  X[x].push_back(y); Y[y].push_back(y);
 int solve(){
  int cnt = 0;
  for(int i=0;i<n;i++){</pre>
   walked.reset();
   if(dfs(i)) cnt++;
  // return how many pair matched
  return cnt;
```

qu[qr++] = s;

vr[s] = true;

for (int i=0; i<n; i++)</pre>

```
| };
                                                                for (int j=0; j<n; j++)</pre>
                                                                 edge[i][j] = 0;
4.3 General Graph Matching
namespace matching {
                                                              void set_edge(int u, int v, int w) {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
                                                               edge[u][v] = edge[v][u] = w;
vector<int> g[kN];
queue<int> q;
                                                              bool SPFA(int u){
                                                               if (onstk[u]) return true;
void Init(int n) {
                                                               stk.PB(u);
 for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
 for (int i = 0; i < n; ++i) g[i].clear();</pre>
                                                               onstk[u] = 1;
                                                               for (int v=0; v<n; v++){</pre>
                                                                if (u != v && match[u] != v && !onstk[v]){
void AddEdge(int u, int v) {
                                                                 int m = match[v]
 g[u].push_back(v);
 g[v].push_back(u);
                                                                 if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
                                                                  dis[m] = dis[u] - edge[v][m] + edge[u][v];
int Find(int u) {
                                                                  onstk[v] = 1;
 return u == fa[u] ? u : fa[u] = Find(fa[u]);
                                                                  stk.PB(v)
                                                                  if (SPFA(m)) return true;
                                                                  stk.pop_back();
int LCA(int x, int y, int n) {
 static int tk = 0; tk++;
                                                                  onstk[v] = 0;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
                                                               onstk[u] = 0;
   v[x] = tk;
                                                               stk.pop_back();
                                                               return false;
   x = Find(pre[match[x]]);
                                                              int solve() {
                                                               // find a match
void Blossom(int x, int y, int 1) {
 while (Find(x) != 1) {
                                                               for (int i=0; i<n; i+=2){
  pre[x] = y, y = match[x];
                                                                match[i] = i+1;
   if (s[y] == 1) q.push(y), s[y] = 0;
                                                                match[i+1] = i;
  if (fa[x] == x) fa[x] = 1;
  if (fa[y] == y) fa[y] = 1;
                                                               while (true){
                                                                int found = 0;
  x = pre[y];
                                                                for (int i=0; i<n; i++)</pre>
                                                                 dis[i] = onstk[i] = 0;
bool Bfs(int r, int n) {
                                                                for (int i=0; i<n; i++){</pre>
 for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;
                                                                 stk.clear()
 while (!q.empty()) q.pop();
                                                                 if (!onstk[i] && SPFA(i)){
 q.push(r);
                                                                  found = 1;
                                                                  while (SZ(stk)>=2){
 s[r] = 0;
 while (!q.empty()) {
                                                                   int u = stk.back(); stk.pop_back();
                                                                   int v = stk.back(); stk.pop_back();
  int x = q.front(); q.pop();
   for (int u : g[x]) {
                                                                   match[u] = v;
   if (s[u] == -1)
                                                                   match[v] = u;
    pre[u] = x, s[u] = 1;
                                                                  }
     if (match[u] == n) {
      for (int a = u, b = x, last; b != n; a = last, b =
                                                                if (!found) break;
      pre[a])
       last = match[b], match[b] = a, match[a] = b;
                                                               int ret = 0;
      return true;
                                                               for (int i=0; i<n; i++)
                                                                ret += edge[i][match[i]];
    q.push(match[u]);
    s[match[u]] = 0;
                                                               return ret>>1;
    } else if (!s[u] && Find(u) != Find(x)) {
    int 1 = LCA(u, x, n);
Blossom(x, u, 1);
                                                             } graph;
                                                             4.5 Minimum Cost Circulation
     Blossom(u, x, 1);
                                                             struct Edge { int to, cap, rev, cost; };
  }
                                                             vector<Edge> g[kN];
                                                             int dist[kN], pv[kN], ed[kN];
 return false;
                                                             bool mark[kN];
                                                             int NegativeCycle(int n) {
int Solve(int n) {
                                                              memset(mark, false, sizeof(mark));
 int res = 0;
                                                              memset(dist, 0, sizeof(dist));
 for (int x = 0; x < n; ++x) {
                                                              int upd = -1;
                                                              for (int i = 0; i <= n; ++i) {</pre>
  if (match[x] == n) res += Bfs(x, n);
                                                               for (int j = 0; j < n; ++j) {
 return res;
                                                                int idx = 0;
                                                                for (auto &e : g[j]) {
}}
                                                                 if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
      Minimum Weight Matching (Clique version)
                                                                  dist[e.to] = dist[j] + e.cost;
struct Graph {
                                                                  pv[e.to] = j, ed[e.to] = idx;
 // 0-base (Perfect Match)
                                                                  if (i == n) {
                                                                   upd = j;
 int n, edge[MXN][MXN]
 int match[MXN], dis[MXN], onstk[MXN];
                                                                   while(!mark[upd])mark[upd]=1,upd=pv[upd];
 vector<int> stk;
                                                                   return upd;
 void init(int _n) {
  n = _n;
```

idx++;

```
}
return -1;
int Solve(int n) {
int rt = -1, ans = 0;
while ((rt = NegativeCycle(n)) >= 0) {
 memset(mark, false, sizeof(mark));
 vector<pair<int, int>> cyc;
while (!mark[rt]) {
   cyc.emplace_back(pv[rt], ed[rt]);
  mark[rt] = true;
   rt = pv[rt];
  reverse(cyc.begin(), cyc.end());
  int cap = kInf;
  for (auto &i : cyc)
  auto &e = g[i.first][i.second];
   cap = min(cap, e.cap);
  for (auto &i : cyc) {
   auto &e = g[i.first][i.second];
   e.cap -= cap;
   g[e.to][e.rev].cap += cap;
   ans += e.cost * cap;
  }
return ans:
```

4.6 Flow Models

- · Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}$.
 - 2. For each edge (x,y,l,u), connect x o y with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect v o T with capacity -in(v).
 - To maximize, connect $t\to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T.If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the
 - maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect t o s with capacity ∞ and let the flow from S to T be f'. If $f+f'\neq \sum_{v\in V, in(v)>0}in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e \mbox{,}$ where f_e corresponds to the flow of edge e on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipartite graph(X, Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise. 2. DFS from unmatched vertices in X.

 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect $x\to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y\to x$ with (cost,cap)=(-c,1)
 - 3. For each edge with c<0, sum these cost as K, then increase d(y)
 - by 1, decrease $d(\boldsymbol{x})$ by 1
 - 4. For each vertex v with d(v) > 0, connect $S \to v$ with (cost, cap) =(0, d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =(0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source s o v, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity
 - 5. For $v \in \mathit{G}$, connect it with sink $v \rightarrow t$ with capacity K + 2T - $(\sum_{e \in E(v)} w(e)) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u,v).
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to \bar{v} .
 - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$. 2. Create edge (u,v) with capacity w with w being the cost of choosing
 - u without choosing v.

- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x, t) with capacity c_x and create edge (s, y) with ca-
- 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4.7 Dinic

```
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
     int to, rev;
     Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS() {
     lv.assign(n, -1);
     queue<int> bfs;
     bfs.push(st); lv[st] = 0;
     while (not bfs.empty()){
       int u = bfs.front(); bfs.pop();
       for (auto e: G[u]) {
         if (e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
         bfs.push(e.to); lv[e.to] = lv[u] + 1;
       }
     return lv[ed] != -1;
  Cap DFS(int u, Cap f){
  if (u == ed) return f;
     Cap ret = 0;
     for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
       auto &e = G[u][i];
       if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
       Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
       if (f == 0) return ret;
     if (ret == 0) lv[u] = -1;
     return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
     G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
  st = st_, ed = ed_; Cap ret = 0;
     while (BFS()) {
       idx.assign(n, 0);
       Cap f = DFS(st, numeric_limits<Cap>::max());
       ret += f;
       if (f == 0) break;
     return ret;
```

Minimum Cost Maximum Flow 4.8

```
class MiniCostMaxiFlow{
 using Cap = int; using Wei = int64_t;
 using PCW = pair<Cap,Wei>;
 static constexpr Cap INF_CAP = 1 << 30;</pre>
 static constexpr Wei INF_WEI = 1LL<<60;</pre>
private:
 struct Edge{
  int to, back;
  Cap cap; Wei wei;
  Edge() {}
  Edge(int a,int b, Cap c, Wei d):
   to(a),back(b),cap(c),wei(d)
```

```
{}
                                                                    const int maxn = 500 + 5;
 };
                                                                    int w[maxn][maxn], g[maxn];
int ori, edd;
                                                                    bool v[maxn], del[maxn];
                                                                    void add_edge(int x, int y, int c) {
vector<vector<Edge>> G;
vector<int> fa, wh;
                                                                    w[x][y] += c; w[y][x] += c;
vector<bool> inq;
 vector<Wei> dis;
                                                                    pair<int, int> phase(int n) {
PCW SPFA(){
                                                                     memset(v, false, sizeof(v));
                                                                     memset(g, 0, sizeof(g));
int s = -1, t = -1;
  fill(inq.begin(),inq.end(),false);
  fill(dis.begin(), dis.end(), INF_WEI);
                                                                     while (true) {
  queue<int> qq; qq.push(ori);
  dis[ori] = 0;
                                                                      int c = -1;
                                                                      for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
  while(not qq.empty()){
   int u=qq.front();qq.pop();
   inq[u] = false;
                                                                        if (c == -1 \mid | g[i] > g[c]) c = i;
   for(int i=0;i<SZ(G[u]);++i){</pre>
                                                                      if (c == -1) break;
    Edge e=G[u][i];
                                                                      v[s = t, t = c] = true;
    int v=e.to; Wei d=e.wei;
                                                                      for (int i = 0; i < n; ++i) {
    if(e.cap <= 0 | |dis[v] <= dis[u] + d)
                                                                       if (del[i] || v[i]) continue;
    dis[v] = dis[u] + d;
                                                                       g[i] += w[c][i];
    fa[v] = u, wh[v] = i;
                                                                      }
    if (inq[v]) continue;
                                                                     return make_pair(s, t);
    qq.push(v);
    inq[v] = true;
                                                                    int mincut(int n) {
                                                                     int cut = 1e9;
  if(dis[edd]==INF_WEI) return {-1, -1};
                                                                     memset(del, false, sizeof(del));
  Cap mw=INF_CAP;
                                                                     for (int i = 0; i < n - 1; ++i) {
  for(int i=edd;i!=ori;i=fa[i])
                                                                      int s, t; tie(s, t) = phase(n);
   mw=min(mw,G[fa[i]][wh[i]].cap);
                                                                      del[t] = true; cut = min(cut, g[t]);
                                                                      for (int j = 0; j < n; ++j) {
  for (int i=edd;i!=ori;i=fa[i]){
                                                                       w[s][j] += w[t][j]; w[j][s] += w[j][t];
   auto &eg=G[fa[i]][wh[i]];
   eg.cap -= mw;
   G[eg.to][eg.back].cap+=mw;
                                                                     return cut;
  return {mw, dis[edd]};
public:
                                                                    5
                                                                          Math
void init(int a,int b,int n){
                                                                         Prime Table
  ori=a,edd=b;
                                                                    1002939109, 1020288887, 1028798297, 1038684299,\\
  G.clear();G.resize(n);
                                                                    1041211027, 1051762951, 1058585963, 1063020809,
  fa.resize(n);wh.resize(n);
                                                                    1147930723, 1172520109, 1183835981, 1187659051,
  inq.resize(n); dis.resize(n);
                                                                    1241251303, 1247184097, 1255940849, 1272759031,
                                                                    1287027493, 1288511629, 1294632499, 1312650799,
 void add_edge(int st, int ed, Cap c, Wei w){
                                                                    1868732623, 1884198443, 1884616807, 1885059541\\
  G[st].emplace_back(ed,SZ(G[ed]),c,w);
                                                                    \begin{array}{c} 1909942399, 1914471137, 1923951707, 1925453197, \\ 1979612177, 1980446837, 1989761941, 2007826547, \\ 2008033571, 2011186739, 2039465081, 2039728567, \end{array}
  G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
                                                                    2093735719, 2116097521, 2123852629, 2140170259,
PCW solve(){
                                                                    3148478261, 3153064147, 3176351071, 3187523093,\\
  Cap cc=0; Wei ww=0;
                                                                    3196772239, 3201312913, 3203063977, 3204840059
  while(true){
                                                                    3210224309, 3213032591, 3217689851, 3218469083, 3219857533, 3231880427, 3235951699, 3273767923,
  PCW ret=SPFA();
                                                                    3276188869, 3277183181, 3282463507, 3285553889,
   if(ret.first==-1) break;
                                                                    3319309027, 3327005333, 3327574903, 3341387953,
   cc+=ret.first:
                                                                    3373293941, 3380077549, 3380892997, 3381118801\\
   ww+=ret.first * ret.second:
                                                                          \lfloor \frac{n}{i} \rfloor Enumeration
                                                                    T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor
  return {cc,ww};
}
                                                                    5.3 ax+by=gcd
} mcmf;
                                                                    // ax+ny = 1, ax+ny == ax == 1 \pmod{n}
                                                                    void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
4.9
      GomoryHu Tree
                                                                     if (y == 0) g=x,a=1,b=0;
int g[maxn];
                                                                     else exgcd(y,x%y,g,b,a),b=(x/y)*a;
vector<edge> GomoryHu(int n){
vector<edge> rt;
                                                                    5.4
                                                                           Pollard Rho
 for(int i=1;i<=n;++i)g[i]=1;</pre>
for(int i=2;i<=n;++i){</pre>
                                                                    // does not work when n is prime
  int t=g[i]
                                                                    // return any non-trivial factor
                  // clear flows on all edge
  flow.reset();
                                                                    llu pollard_rho(llu n){
  rt.push_back({i,t,flow(i,t)});
                                                                     static auto f=[](llu x,llu k,llu m){
  flow.walk(i); // bfs points that connected to i (use
                                                                      return add(k,mul(x,x,m),m);
    edges not fully flow)
  for(int j=i+1;j<=n;++j){</pre>
                                                                     if (!(n&1)) return 2;
   if(g[j]==t && flow.connect(j))g[j]=i; // check if i
                                                                     mt19937 rnd(120821011);
    can reach j
                                                                     while(true){
  }
                                                                      llu y=2, yy=y, x=rnd()%n, t=1;
                                                                      for(llu sz=2;t==1;sz<<=1) {</pre>
return rt;
                                                                       for(llu i=0;i<sz;++i){</pre>
                                                                         if(t!=1)break;
```

yy=f(yy,x,n);

t=gcd(yy>y?yy-y:y-yy,n);

4.10 Global Min-Cut

```
}
y=yy;
}
if(t!=1&&t!=n) return t;
}
}
```

5.5 Pi Count (Linear Sieve)

```
static constexpr int N = 1000000 + 5;
11d pi[N];
vector<int> primes;
bool sieved[N];
11d cube_root(11d x){
 1ld s=cbrt(x-static_cast<long double>(0.1));
 while(s*s*s <= x) ++s;
 return s-1;
11d square_root(11d x){
lld s=sqrt(x-static_cast<long double>(0.1));
 while(s*s \ll x) ++s;
 return s-1;
void init(){
 primes.reserve(N)
 primes.push_back(1);
 for(int i=2;i<N;i++) {</pre>
  if(!sieved[i]) primes.push_back(i);
  pi[i] = !sieved[i] + pi[i-1];
  for(int p: primes) if(p > 1) {
   if(p * i >= N) break;
   sieved[p * i] = true;
}
   if(p % i == 0) break;
  }
11d phi(11d m, 11d n) {
static constexpr int MM = 80000, NN = 500;
static lld val[MM][NN];
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
 lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
 if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
 return ret;
11d pi_count(11d);
11d P2(11d m, 11d n) {
 11d sm = square_root(m), ret = 0;
 for(lld i = n+1;primes[i]<=sm;i++)</pre>
  ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
 return ret;
11d pi_count(11d m) {
 if(m < N) return pi[m];</pre>
 11d n = pi_count(cube_root(m));
return phi(m, n) + n - 1 - P2(m, n);
```

5.6 Strling Number

5.6.1 First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1)\dots(x+n-1) = \sum_{k=0}^n S_1(n,k)x^k$$

$$g(x) = x(x+1)\dots(x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.6.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

```
S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}
```

5.7 Range Sieve

```
const int MAX_SQRT_B = 50000;
const int MAX_L = 200000 + 5;

bool is_prime_small[MAX_SQRT_B];
bool is_prime[MAX_L];

void sieve(lld l, lld r){
   // [l, r)
   for(lld i=2;i*i<r;i++) is_prime_small[i] = true;
   for(lld i=1;i<r;i++) is_prime[i-1] = true;
   if(l=1) is_prime[0] = false;
   for(lld i=2;i*i<r;i++){
    if(!is_prime_small[i]) continue;
   for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;
   for(lld j=std::max(2LL, (l+i-1)/i)*i;j<r;j+=i)
        is_prime[j-l]=false;
   }
}</pre>
```

5.8 Miller Rabin

```
bool isprime(llu x){
 static llu magic[]={2,325,9375,28178,\
          450775,9780504,1795265022};
 static auto witn=[](llu a,llu u,llu n,int t)
 ->bool{
  if (!(a = mpow(a%n,u,n)))return 0;
  while(t--){
   1lu a2=mul(a,a,n);
   if(a2==1 && a!=1 && a!=n-1)
    return 1;
   a = a2;
  }
  return a!=1;
 if(x<2)return 0;</pre>
 if(!(x&1))return x==2;
 llu x1=x-1; int t=0;
 while(!(x1&1))x1>>=1,t++;
 for(llu m:magic)if(witn(m,x1,x,t))return 0;
 return 1:
```

5.9 Extended Euler

```
a^b \equiv \begin{cases} a^b \mod \varphi(m) + \varphi(m) & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m
```

5.10 Gauss Elimination

```
void gauss(vector<vector<double>> &d) {
  int n = d.size(), m = d[0].size();
  for (int i = 0; i < m; ++i) {
    int p = -1;
    for (int j = i; j < n; ++j) {
      if (fabs(d[j][i]) < eps) continue;
      if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
    }
    if (p == -1) continue;
    for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
    for (int j = 0; j < n; ++j) {
      if (i == j) continue;
      double z = d[j][i] / d[i][i];
      for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
    }
}</pre>
```

5.11 Fast Fourier Transform

```
const int mod = 10000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
   static_assert (M1 <= M2 && M2 <= M3);
   constexpr int64_t r12 = modpow(M1, M2-2, M2);</pre>
```

```
constexpr int64_t r13 = modpow(M1, M3-2, M3);
                                                                 cplx b2 = (fb[i] - fb[j].conj()) * r4;
  constexpr int64_t r23 = modpow(M2, M3-2, M3);
                                                                 if (i != j) {
                                                                  cplx c1 = (fa[j] + fa[i].conj());
  constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
                                                                  cplx c2 = (fa[j] - fa[i].conj()) * r2;
cplx d1 = (fb[j] + fb[i].conj()) * r3;
  B = (B - A + M2) * r12 % M2;
 C = (C - A + M3) * r13 % M3;
 C = (C - B + M3) * r23 % M3
                                                                  cplx d2 = (fb[j] - fb[i].conj()) * r4;
  return (A + B * M1 + C * M1M2) % mod;
                                                                  fa[i] = c1 * d1 + c2 * d2 * r5;
                                                                  fb[i] = c1 * d2 + c2 * d1:
namespace fft {
                                                                 fa[j] = a1 * b1 + a2 * b2 * r5;
                                                                 fb[j] = a1 * b2 + a2 * b1;
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
                                                                fft(fa, sz), fft(fb, sz);
cplx omega[maxn + 1];
                                                                vector<int> res(sz);
                                                                for (int i = 0; i < sz; ++i) {
void prefft() {
for (int i = 0; i <= maxn; i++)</pre>
                                                                 long long a = round(fa[i].re), b = round(fb[i].re),
  omega[i] = cplx(cos(2 * pi * j / maxn),
                                                                       c = round(fa[i].im);
                                                                 res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
     sin(2 * pi * j / maxn));
                                                                }
void fft(vector<cplx> &v, int n) {
                                                                return res;
 int z = __builtin_ctz(n) - 1;
                                                               }}
 for (int i = 0; i < n; ++i) {
                                                               5.12 Chinese Remainder
  int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^=(i >> j & 1)<<(z - j);
                                                               lld crt(lld ans[], lld pri[], int n){
  if (x > i) swap(v[x], v[i]);
                                                                lld M = 1, ret = 0;
                                                                for(int i=0;i<n;i++) M *= pri[i];</pre>
 for (int s = 2; s <= n; s <<= 1) {
                                                                for(int i=0;i<n;i++)</pre>
  int z = s >> 1;
                                                                 lld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
  for (int i = 0; i < n; i += s) {
                                                                 ret += (ans[i]*(M/pri[i])%M * iv)%M;
   for (int k = 0; k < z; ++k) {
  cplx x = v[i + z + k] * omega[maxn / s * k];
                                                                 ret %= M;
    v[i + z + k] = v[i + k] - x;
                                                                return ret;
    v[i + k] = v[i + k] + x;
                                                               }
                                                               /*
                                                               Another:
                                                               x = a1 \% m1
                                                               x = a2 \% m2
void ifft(vector<cplx> &v, int n) {
                                                               g = gcd(m1, m2)
 fft(v, n);
                                                               assert((a1-a2)%g==0)
 reverse(v.begin() + 1, v.end());
                                                               [p, q] = exgcd(m2/g, m1/g)
 for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);
                                                               return a2+m2*(p*(a1-a2)/g)
                                                               \theta \ll x \ll 1cm(m1, m2)
VL convolution(const VI &a, const VI &b) {
                                                               */
 // Should be able to handle N <= 10^5, C <= 10^4
 int sz = 1;
                                                                      Berlekamp Massey
 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
                                                               // x: 1-base, p[]: 0-base
 vector<cplx> v(sz);
 for (int i = 0; i < sz; ++i) {
                                                               template<size t N>
                                                               vector<llf> BM(llf x[N], size_t n){
  double re = i < a.size() ? a[i] : 0;</pre>
                                                                size_t f[N]={0},t=0;11f d[N];
  double im = i < b.size() ? b[i] : 0;</pre>
                                                                vector<llf> p[N];
  v[i] = cplx(re, im);
                                                                for(size_t i=1,b=0;i<=n;++i) {</pre>
                                                                 for(size_t j=0;j<p[t].size();++j)
d[i]+=x[i-j-1]*p[t][j];</pre>
 fft(v, sz);
 for (int i = 0; i <= sz / 2; ++i) {
                                                                 if(abs(d[i]-=x[i])<=EPS)continue;</pre>
  int j = (sz - i) & (sz - 1);
                                                                 f[t]=i;if(!t){p[++t].resize(i);continue;}
 cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
                                                                 vector<llf> cur(i-f[b]-1);
    * cplx(0, -0.25);
                                                                 11f k=-d[i]/d[f[b]];cur.PB(-k);
  if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
                                                                 for(size_t j=0;j<p[b].size();j++)
  cur.PB(p[b][j]*k);</pre>
    ].conj()) * cplx(\hat{0}, -0.25);
  v[i] = x;
                                                                 if(cur.size()<p[t].size())cur.resize(p[t].size());</pre>
                                                                 for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];</pre>
 ifft(v, sz);
                                                                 if(i-f[b]+p[b].size()>=p[t].size()) b=t;
 VL c(sz);
                                                                 p[++t]=cur;
 for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
 return c;
                                                                return p[t];
VI convolution_mod(const VI &a, const VI &b, int p) {
                                                               5.14 NTT
 while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
 vector<cplx> fa(sz), fb(sz);
                                                               template <int mod, int G, int maxn>
 for (int i = 0; i < (int)a.size(); ++i)</pre>
                                                               struct NTT {
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
                                                                static_assert (maxn == (maxn & -maxn));
 for (int i = 0; i < (int)b.size(); ++i)</pre>
                                                                int roots[maxn];
  fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
                                                                NTT () {
 fft(fa, sz), fft(fb, sz);
                                                                 int r = modpow(G, (mod - 1) / maxn);
 double r = 0.25 / sz;
                                                                 for (int i = maxn >> 1; i; i >>= 1) {
 cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
                                                                  roots[i] = 1;
                                                                  for (int j = 1; j < i; j++)
  int j = (sz - i) & (sz - 1);
                                                                   roots[i + j] = modmul(roots[i + j - 1], r);
 cplx a1 = (fa[i] + fa[j].conj());
cplx a2 = (fa[i] - fa[j].conj()) * r2;
                                                                  r = modmul(r, r);
  cplx b1 = (fb[i] + fb[j].conj()) * r3;
```

```
// n must be 2^k, and 0 \le F[i] < mod
                                                                 Poly Sqrt() const { // Jacobi(coef[0], P) = 1
 void inplace_ntt(int n, int F[], bool inv = false) {
                                                                  if (n()==1) return {QuadraticResidue(coef[0], P)};
  for (int i = 0, j = 0; i < n; i++) {
  if (i < j) swap(F[i], F[j]);</pre>
                                                                  Poly X = Poly(*this, (n()+1) / 2).Sqrt().isz(n());
                                                                  return X.iadd(Mul(X.Inv()).isz(n())).imul(P/2+1);
   for (int k = n > 1; (j^k < k; k > = 1);
                                                                 pair<Poly, Poly> DivMod(const Poly &rhs) const {
                                                                  // (rhs.)back() != 0
  for (int s = 1; s < n; s *= 2) {
   for (int i = 0; i < n; i += s * 2) {
                                                                  if (n() < rhs.n()) return {{0}, *this};</pre>
    for (int j = 0; j < s; j++) {
                                                                  const int _n = n() - rhs.n() + 1;
     int a = F[i+j];
                                                                  Poly X(rhs); X.irev().isz(_n);
     int b = modmul(F[i+j+s], roots[s+j]);
                                                                  Poly Y(*this); Y.irev().isz(_n);
     F[i+j] = modadd(a, b); // a + b
                                                                  Poly Q = Y.Mul(X.Inv()).isz(_n).irev();
                                                                  X = rhs.Mul(Q), Y = *this;
fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;</pre>
     F[i+j+s] = modsub(a, b); // a - b
                                                                  return {Q, Y.isz(max(1, rhs.n() - 1))};
   }
  if (inv) {
                                                                 Poly Dx() const {
                                                                 Poly ret(n() - 1);
   int invn = modinv(n);
                                                                  fi(0, ret.n()) ret[i] = (i + 1) * coef[i + 1] % P;
   for (int i = 0; i < n; i++)</pre>
    F[i] = modmul(F[i], invn);
                                                                  return ret.isz(max(1, ret.n()));
   reverse(F + 1, F + n);
                                                                 Poly Sx() const {
  Poly ret(n() + 1);
                                                                  fi(0, n()) ret[i + 1]=ntt.minv(i + 1)*coef[i] % P;
};
const int P=2013265921, root=31;
                                                                  return ret;
const int MAXN=1<<20;</pre>
NTT<P, root, MAXN> ntt;
                                                                 Poly _tmul(int nn, const Poly &rhs) const {
                                                                  Poly Y = Mul(rhs).isz(n() + nn - 1);
5.15 Polynomial Operations
                                                                  return Poly(Y.data() + n() - 1, nn);
using VL = vector<LL>
#define fi(s, n) for (int i=int(s); i<int(n); ++i)</pre>
                                                                 VL _eval(const VL &x, const auto up)const{
#define Fi(s, n) for (int i=int(n); i>int(s); --i)
                                                                  const int _n = (int)x.size();
int n2k(int n) {
                                                                  if (!_n) return {};
                                                                  vector<Poly> down(_n * 2);
int sz = 1; while (sz < n) sz <<= 1;</pre>
 return sz;
                                                                  down[1] = DivMod(up[1]).second;
                                                                  fi(2,_n*2) down[i]=down[i/2].DivMod(up[i]).second;
                                                                  /* down[1] = Poly(up[1]).irev().isz(n()).Inv().irev()
template<int MAXN, LL P, LL RT> // MAXN = 2^k
                                                                      _tmul(_n, *this)
struct Poly { // coefficients in [0, P)
 static NTT<MAXN, P, RT> ntt;
                                                                  fi(2, _n * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
                                                                    1, down[i / 2]); */
 VL coef;
 int n() const { return coef.size(); } // n()>=1
                                                                  VL y(_n);
                                                                  fi(0, _n) y[i] = down[_n + i][0];
 LL *data() { return coef.data(); }
 const LL *data() const { return coef.data(); }
                                                                  return y;
 LL &operator[](size_t i) { return coef[i]; }
                                                                 static vector<Poly> _tree1(const VL &x) {
 const LL &operator[](size_t i)const{return coef[i];}
                                                                  const int _n = (int)x.size();
 Poly(initializer_list<LL> a) : coef(a) { }
                                                                  vector<Poly> up(_n * 2);
 explicit Poly(int _n = 1) : coef(_n) { }
Poly(const LL *arr, int _n) : coef(arr, arr + _n) {}
Poly(const Poly &p, int _n) : coef(_n) {
                                                                  fi(0, _n) up[_n + i] = \{(x[i] ? P - x[i] : 0), 1\};
                                                                  Fi(0, _n-1) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
  copy_n(p.data(), min(p.n(), _n), data());
                                                                  return up;
 Poly& irev(){return reverse(data(),data()+n()),*this;}
                                                                 VL Eval(const VL&x)const{return _eval(x,_tree1(x));}
                                                                 static Poly Interpolate(const VL &x, const VL &y) {
 Poly& isz(int _n) { return coef.resize(_n), *this; }
 Poly& iadd(const Poly &rhs) { // n() == rhs.n()
                                                                  const int _n = (int)x.size();
                                                                  vector<Poly> up = _tree1(x), down(_n * 2);
VL z = up[1].Dx()._eval(x, up);
  fi(0, n()) if ((coef[i]+=rhs[i]) >= P)coef[i]-=P;
  return *this;
                                                                  fi(0, _n) z[i] = y[i] * ntt.minv(z[i]) % P;
                                                                  fi(0, _n) down[_n + i] = \{z[i]\};

Fi(0, _n-1) down[i] = down[i * 2].Mul(up[i * 2 + 1])
 Poly& imul(LL k) {
 fi(0, n()) coef[i] = coef[i] * k % P;
                                                                   .iadd(down[i * 2 + 1].Mul(up[i * 2]));
  return *this;
                                                                  return down[1];
 Poly Mul(const Poly &rhs) const {
                                                                 Poly Ln() const { // coef[0] == 1
  const int _n = n2k(n() + rhs.n() - 1);
 Poly X(*this, _n), Y(rhs, _n);
ntt(X.data(), _n), ntt(Y.data(),
fi(0, _n) X[i] = X[i] * Y[i] % P;
                                                                  return Dx().Mul(Inv()).Sx().isz(n());
                                                                 Poly Exp() const \{ // coef[0] == 0 \}
                                                                  if (n() == 1) return {1};
  ntt(X.data(), _n, true);
                                                                  Poly X = Poly(*this, (n() + 1)/2).Exp().isz(n());
Poly Y = X.Ln(); Y[0] = P - 1;
  return X.isz(n() + rhs.n() - 1);
                                                                  fi(0, n()) if((Y[i] = coef[i] - Y[i]) < 0)Y[i]+=P;
 Poly Inv() const { // coef[0] != 0
                                                                  return X.Mul(Y).isz(n());
 if (n() == 1) return {ntt.minv(coef[0])};
  const int _n = n2k(n() * 2);
  Poly Xi = Poly(*this, (n() + 1)/2).Inv().isz(_n);
                                                                 Poly Pow(const string &K) const {
  Poly Y(*this, _n);
                                                                  int nz = 0;
                                                                  while (nz < n() && !coef[nz]) ++nz;</pre>
  ntt(Xi.data(), _n), ntt(Y.data(), _n);
  fi(0, _n) {
Xi[i] *= (2 - Xi[i] * Y[i]) % P
                                                                  LL nk = 0, nk2 = 0;
                                                                  for (char c : K) {
                                                                  nk = (nk * 10 + c - '0') % P;
nk2 = nk2 * 10 + c - '0';
   if((Xi[i] \% = P) < 0) Xi[i] += P;
                                                                   if (nk2 * nz >= n()) return Poly(n());
  ntt(Xi.data(), _n, true);
  return Xi.isz(n());
                                                                   nk2 %= P - 1;
```

```
if (!nk && !nk2) return Poly({1}, n());
                                                                    if (y % g != 0) return -1;
  Poly X(data() + nz, n() - nz * nk2);
                                                                    t /= g, y /= g, M /= g;
  LL \times 0 = X[0]
                                                                    Int h = 0, gs = 1;
  return X.imul(ntt.minv(x0)).Ln().imul(nk).Exp()
                                                                    for (; h * h < M; ++h) gs = gs * x % M;
                                                                    unordered_map<Int, Int> bs;
   .imul(ntt.mpow(x0, nk2)).irev().isz(n()).irev();
                                                                    for (Int s = 0; s < h; bs[y] = ++s)
 Poly InvMod(int L) { // (to evaluate linear recursion)
                                                                     y = y * x % M;
 Poly R\{1, 0\}; // *this * R mod x^L = 1 (*this[0] ==
                                                                    for (Int s = 0; s < M; s += h) {
                                                                      t = t * gs % M;
  for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                                      if (bs.count(t)) return c + s + h - bs[t];
   Poly 0 = R.Mul(Poly(data(), min(2 << level, n())));
   Poly Q(2 \ll level); Q[0] = 1;
                                                                    return -1;
   for (int j = (1 << level); j < (2 << level); ++j)
Q[j] = (P - O[j]) % P;</pre>
   R = R.Mul(Q).isz(4 << level);
                                                                 5.18
                                                                        FloorSum
  }
                                                                 // @param n `n < 2^32`
  return R.isz(L);
                                                                 // @param m `1 <= m < 2^32`
 \textbf{static} \  \, \textbf{LL} \  \, \textbf{LinearRecursion}(\textbf{const} \  \, \textbf{VL\&a}, \textbf{const} \  \, \textbf{VL\&c}, \textbf{LL} \  \, \textbf{n}) \, \{
                                                                  // @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
                                                                 llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
  // a_n = \sum_{j=0}^{n-j} a_{n-j}
  const int k = (int)a.size();
                                                                   llu ans = 0:
                                                                   while (true)
  assert((int)c.size() == k + 1);
                                                                   if (a >= m) {
  Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
                                                                     ans += n * (n - 1) / 2 * (a / m); a %= m;
  fi(1, k + 1) C[k - i] = c[i] ? P - c[i] : 0;
  C[k] = 1
                                                                    if (b >= m) {
  while (n) {
                                                                     ans += n * (b / m); b %= m;
   if (n % 2) W = W.Mul(M).DivMod(C).second;
   n /= 2, M = M.Mul(M).DivMod(C).second;
                                                                    llu y_max = a * n + b;
                                                                    if (y_max < m) break;</pre>
  LL ret = 0:
                                                                    // y_max < m * (n + 1)
  fi(0, k) ret = (ret + W[i] * a[i]) % P;
                                                                    // floor(y_max / m) <= n
  return ret;
                                                                   n = (11u)(y_max / m), b = (11u)(y_max % m);
                                                                    swap(m, a);
}:
#undef fi
                                                                   return ans;
#undef Fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
                                                                 11d floor_sum(11d n, 11d m, 11d a, 11d b) {
template<> decltype(Poly_t::ntt) Poly_t::ntt = {};
                                                                  11u ans = 0;
5.16 FWT
                                                                   if (a < 0) {
                                                                   11u \ a2 = (a \% m + m) \% m;
/* xor convolution:
* x = (x0, x1) , y = (y0, y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                   ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
* x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1)
* z = (1/2) * z''
                                                                   if (b < 0) {
                                                                   11u b2 = (b \% m + m) \% m;
                                                                    ans -= 1ULL * n * ((b2 - b) / m);
 * or convolution:
                                                                    b = b2:
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                                   return ans + floor_sum_unsigned(n, m, a, b);
 * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
const LL MOD = 1e9+7;
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
                                                                 5.19 Quadratic residue
 for( int d = 1 ; d < N ; d <<= 1 ) {
                                                                 struct S {
  int d2 = d << 1;
                                                                  int MOD, w;
  for( int s = 0 ; s < N ; s += d2 )
                                                                   int64_t x, y;
   for( int i = s , j = s+d ; i < s+d ; i++, j++ ){ LL ta = x[i] , tb = x[j];
                                                                   S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
                                                                    : MOD(m), w(w_), x(x_), y(y_) {}
    x[ i ] = ta+tb;
                                                                   S operator*(const S &rhs) const {
    x[ j ] = ta-tb;
    if( x[ i ] >= MOD ) x[ i ] -= MOD;
if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
                                                                   int w_{-} = w
                                                                    if (w<sub>_</sub> == -1) w<sub>_</sub> = rhs.w;
                                                                    assert(w_ != -1 and w_ == rhs.w);
                                                                   return { MOD, w_,
(x * rhs.x + y * rhs.y % MOD * w) % MOD,
 }
 if( inv )
                                                                     (x * rhs.y + y * rhs.x) % MOD };
  for( int i = 0 ; i < N ; i++ ) {</pre>
   x[`i ] *= inv(`N, MOD );
x[ i ] %= MOD;
                                                                 int get_root(int n, int P) {
  }
                                                                   if (P == 2 or n == 0) return n;
                                                                    if (qpow(n, (P - 1) / 2, P) != 1) return -1;
5.17 DiscreteLog
                                                                    auto check = [&](int x) {
                                                                   return qpow(x, (P - 1) / 2, P); };
if (check(n) == P-1) return -1;
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
                                                                    int64_t a; int w; mt19937 rnd(7122);
  // x^? \equiv y (mod M)
  Int t = 1, c = 0, g = 1;
                                                                    do { a = rnd() % P;
                                                                     w = ((a * a - n) & P + P) & P;
  for (Int M_ = M; M_ > 0; M_ >>= 1)
                                                                    } while (check(w) != P - 1);
    g = g * x % M;
                                                                    return qpow(S(P, w, a, 1), (P + 1) / 2).x;
  for (g = gcd(g, M); t % g != 0; ++c) {
    if (t == y) return c;
    t = t * x % M;
                                                                 5.20 De-Bruijn
```

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
  if (n \% p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
 } else {
  aux[t] = aux[t - p];
  db(t + 1, p, n, k);
for (int i = aux[t - p] + 1; i < k; ++i) {
   aux[t] = i;
   db(t + 1, t, n, k);
  }
 }
int de_bruijn(int k, int n) {
  // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
 if (k == 1) {
  res[0] = 0;
  return 1:
 for (int i = 0; i < k * n; i++) aux[i] = 0;
 sz = 0:
 db(1, 1, n, k);
 return sz;
```

5.21 Simplex Construction

Standard form: maximize $\sum_{1 \leq i \leq n} c_i x_i$ such that for all $1 \leq j \leq m$, $\sum_{1 < i < n} A_{ji} x_i \leq b_j$ and $x_i \geq 0$ for all $1 \leq i \leq n$.

- 1. In case of minimization, let $c_i' = -c_i$
- 2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- 3. $\sum_{1 \leq i \leq n} A_{ji} x_i = b_j$
 - $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$
 - $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

5.22 Simplex

```
namespace simplex {
// maximize c^Tx under Ax <= B</pre>
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
double inv = 1.0 / d[r][s];
for (int i = 0; i < m + 2; ++i)
  for (int j = 0; j < n + 2; ++j)
  if (i != r && j != s)
   d[i][j] -= d[r][j] * d[i][s] * inv;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
int x = m + z;
while (true) {
 int s = -1:
  for (int i = 0; i <= n; ++i) {
  if (!z && q[i] == -1) continue;
  if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
  if (d[x][s] > -eps) return true;
  int r = -1
  for (int i = 0; i < m; ++i) {
  if (d[i][s] < eps) continue;
if (r == -1 || \</pre>
   d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
 pivot(r, s);
```

```
VD solve(const VVD &a, const VD &b, const VD &c) {
m = b.size(), n = c.size();
d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)
 for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)</pre>
 p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];
 q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
 for (int i = 1; i < m; ++i)
  if (d[i][n + 1] < d[r][n + 1]) r = i;
 if (d[r][n + 1] < -eps) {
  pivot(r, n)
  if (!phase(1) \mid \mid d[m + 1][n + 1] < -eps)
   return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
   int`s = min_element(d[i].begin(), d[i].end() - 1)
        - d[i].begin();
   pivot(i, s);
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)</pre>
 if (p[i] < n) x[p[i]] = d[i][n + 1];
 return x;
```

5.23 Charateristic Polynomial

```
vector<vector<int>>> Hessenberg(const vector<vector<int</pre>
    >> &A) {
 int N = A.size();
 vector<vector<int>> H = A;
 for (int i = 0; i < N - 2; ++i) {
  if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {
    if (H[j][i]) {
     for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
     ][k])
     for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k]
    ][j]);
     break;
    }
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
   * H[i + 1][k] * (kP - coef)) % kP;
for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
  }
 }
 return H;
vector<int> CharacteristicPoly(const vector<vector<int</pre>
    >> &A) {
 int N = A.size();
 auto H = Hessenberg(A);
 for (int i = 0; i < N; ++i) {
  for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
 vector<vector<int>> P(N + 1, vector<int>(N + 1));
 P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
  P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
    1];
  int val = 1;
  for (int j = i - 1; j >= 0; --j)
   int coef = 1LL * val * H[j][i
                                     _ i] % kP:
   for (int k = 0; k <= j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
```

```
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 if (N & 1) {
                                                                REP(i,n) REP(j,n) flag[i][j] = 0;
  for (int i = 0; i \le N; ++i) P[N][i] = kP - P[N][i];
                                                                vector<Face> now;
                                                                now.emplace_back(0,1,2);
                                                                now.emplace_back(2,1,0);
for (int i=3; i<n; i++){
 return P[N];
}
                                                                 ftop++; vector<Face> next;
                                                                 REP(j, SZ(now)) {
  Face& f=now[j]; int ff = 0;
     Geometry
6
     Basic Geometry
                                                                  ld d=(pt[i]-pt[f.a]).dot(
                                                                     ver(pt[f.a], pt[f.b], pt[f.c]));
using coord_t = int;
                                                                  if (d <= 0) next.push_back(f);</pre>
using Real = double;
                                                                  if (d > 0) ff=ftop;
using Point = std::complex<coord_t>;
                                                                  else if (d < 0) ff=-ftop</pre>
int sgn(coord_t x) {
                                                                  flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
return (x > 0) - (x < 0);
                                                                 REP(j, SZ(now)) {
  Face& f=now[j];
coord_t dot(Point a, Point b) {
return real(conj(a) * b);
                                                                  if (flag[f.a][f.b] > 0 \&\&
                                                                     flag[f.a][f.b] != flag[f.b][f.a])
coord_t cross(Point a, Point b) {
                                                                   next.emplace_back(f.a,f.b,i);
return imag(conj(a) * b);
                                                                  if (flag[f.b][f.c] > 0 &&
                                                                     flag[f.b][f.c] != flag[f.c][f.b])
int ori(Point a, Point b, Point c) {
                                                                   next.emplace_back(f.b,f.c,i);
return sgn(cross(b - a, c - a));
                                                                  if (flag[f.c][f.a] > 0 &&
                                                                    flag[f.c][f.a] != flag[f.a][f.c])
bool operator<(const Point &a, const Point &b) {</pre>
                                                                   next.emplace_back(f.c,f.a,i);
 return real(a) != real(b)
  ? real(a) < real(b) : imag(a) < imag(b);</pre>
                                                                 now=next;
int argCmp(Point a, Point b) {
                                                                return now;
 // -1 / 0 / 1 <-> < / == / > (atan2)
 int qa = (imag(a) == 0
   ? (real(a) < 0 ? 3 : 1) : (imag(a) < 0 ? 0 : 2));
                                                               6.4 2D Farthest Pair
 int qb = (imag(b) == 0
                                                               // stk is from convex hull
   ? (real(b) < 0 ? 3 : 1) : (imag(b) < 0 ? 0 : 2));
 if (qa != qb)
                                                               n = (int)(stk.size());
                                                               int pos = 1, ans = 0; stk.push_back(stk[0]);
  return sgn(qa - qb);
                                                               for(int i=0;i<n;i++) {
  while(abs(cross(stk[i+1]-stk[i],</pre>
 return sgn(cross(b, a));
                                                                  stk[(pos+1)%n]-stk[i])) >
template <typename V> Real area(const V & pt) {
 coord_t ret = 0;
                                                                  abs(cross(stk[i+1]-stk[i],
                                                                 \begin{array}{l} stk[pos] - stk[i]))) \ pos = (pos+1)\%n; \\ ans = max(\{ans, \ dis(stk[i], \ stk[pos]), \end{array} 
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
                                                                 dis(stk[i+1], stk[pos])});
 return ret / 2.0;
                                                               6.5 2D Closest Pair
6.2 2D Convex Hull
                                                               struct cmp_y {
template<typename PT>
vector<PT> buildConvexHull(vector<PT> d) {
                                                                bool operator()(const P& p, const P& q) const {
 sort(ALL(d), [](const PT& a, const PT& b){
                                                                 return p.y < q.y;</pre>
   return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
 vector<PT> s(SZ(d)<<1);
                                                               multiset<P, cmp_y> s;
 int o = 0:
                                                               void solve(P a[], int n) {
 for(auto p: d) {
  while(o>=2 && cross(p-s[o-2], s[o-1]-s[o-2])<=0)
                                                                sort(a, a + n, [](const P& p, const P& q) {
                                                                 return tie(p.x, p.y) < tie(q.x, q.y);</pre>
   0--
  s[o++] = p;
                                                                llf d = INF; int pt = 0;
                                                                for (int i = 0; i < n; ++i) {
 for(int i=SZ(d)-2, t = o+1;i>=0;i--){
  while(o>=t&&cross(d[i]-s[o-2],s[o-1]-s[o-2])<=0)
                                                                 while (pt < i and a[i].x - a[pt].x >= d)
                                                                  s.erase(s.find(a[pt++]));
   0--
                                                                 auto it = s.lower_bound(P(a[i].x, a[i].y - d));
  s[o++] = d[i];
                                                                 while (it != s.end() and it->y - a[i].y < d)</pre>
                                                                  d = min(d, dis(*(it++), a[i]));
 s.resize(o-1);
                                                                 s.insert(a[i]);
 return s;
                                                               }
6.3 3D Convex Hull
                                                                     kD Closest Pair (3D ver.)
                                                               6.6
// return the faces with pt indexes
int flag[MXN][MXN];
                                                               llf solve(vector<P> v) {
                                                                shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,</pre>
struct Point{
 ld x,y,z;
 Point operator * (const ld &b) const {
                                                                 unordered_map<lld, int>>> m;
  return (Point) {x*b, y*b, z*b};}
                                                                llf d = dis(v[0], v[1]);
                                                                auto Idx = [&d] (11f x) -> 11d {
 Point operator * (const Point &b) const {
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
                                                                 return round(x * 2 / d) + 0.1; };
 }
                                                                auto rebuild_m = [&m, &v, &Idx](int k) {
```

m.clear();

}; rebuild_m(2);

for (int i = 0; i < k; ++i)
m[Idx(v[i].x)][Idx(v[i].y)]</pre>

[Idx(v[i].z)] = i;

Point ver(Point a, Point b, Point c) {
return (b - a) * (c - a);}

int n = SZ(pt), ftop = 0;

vector<Face> convex_hull_3D(const vector<Point> pt) {

```
for (size_t i = 2; i < v.size(); ++i) -</pre>
                                                                      && ori(R.st, R.ed, pt.front()) < 0) \
 const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
                                                                      pt.pop_front(), que.pop_front();
     kz = Idx(v[i].z); bool found = false;
                                                                    POP(lines[i], lines[i])
  for (int dx = -2; dx <= 2; ++dx) {
                                                                    pt.push_back(intersect(que.back(), lines[i]));
   const 11d nx = dx + kx;
                                                                    que.push_back(lines[i]);
   if (m.find(nx) == m.end()) continue;
                                                                 POP(que.front(), que.back())
if (que.size() <= 1 ||
   auto& mm = m[nx];
   for (int dy = -2; dy <= 2; ++dy) {
    const 11d ny = dy + ky;
                                                                    argCmp(que.front().dir, que.back().dir) == 0)
    if (mm.find(ny) == mm.end()) continue;
                                                                    return 0;
    auto& mmm = mm[ny];
                                                                  pt.push_back(intersect(que.front(), que.back()));
    for (int dz = -2; dz <= 2; ++dz) {
                                                                  return area(pt);
     const 11d nz = dz + kz;
     if (mmm.find(nz) == mmm.end()) continue;
     const int p = mmm[nz];
                                                               6.9 Minkowski Sum
     if (dis(v[p], v[i]) < d) {
  d = dis(v[p], v[i]);</pre>
                                                               vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
                                                                hull(A), hull(B);
      found = true;
                                                                vector<pll> C(1, A[0] + B[0]), s1, s2;
for(int i = 0; i < SZ(A); ++i)</pre>
     }
                                                                  s1.pb(A[(i + 1) % SZ(A)] - A[i]);
                                                                for(int i = 0; i < SZ(B); i++)
s2.pb(B[(i + 1) % SZ(B)] - B[i]);
  if (found) rebuild_m(i + 1);
                                                                 for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
 else m[kx][ky][kz] = i;
                                                                 if (p2 >= SZ(B)
                                                                    || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
return d;
                                                                   C.pb(C.back() + s1[p1++]);
                                                                  else
                                                                  C.pb(C.back() + s2[p2++]);
      Simulated Annealing
                                                                 return hull(C), C;
11f anneal() {
mt19937 rnd_engine( seed );
uniform_real_distribution< llf > rnd( 0, 1 );
                                                               6.10 Intersection of line and Circle
const 11f dT = 0.001;
                                                               vector<pdd> line_interCircle(const pdd &p1,
 // Argument p
                                                                    const pdd &p2,const pdd &c,const double r){
11f S_cur = calc( p ), S_best = S_cur;
for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
                                                                 pdd ft=foot(p1,p2,c),vec=p2-p1;
                                                                 double dis=abs(c-ft);
 // Modify p to p_prime
                                                                 if(fabs(dis-r)<eps) return vector<pdd>{ft};
 const 11f S_prime = calc( p_prime );
                                                                 if(dis>r) return {};
 const llf delta_c = S_prime - S_cur;
llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
                                                                 vec=vec*sqrt(r*r-dis*dis)/abs(vec);
                                                                 return vector<pdd>{ft+vec,ft-vec};
 if ( rnd( rnd_engine ) <= prob )</pre>
 S_cur = S_prime, p = p_prime;
if ( S_prime < S_best ) // find min</pre>
                                                               6.11 Intersection of Polygon and Circle
  S_best = S_prime, p_best = p_prime;
                                                               // Divides into multiple triangle, and sum up
                                                               // test by HDU2892
return S_best;
                                                               const double PI=acos(-1);
                                                               double _area(pdd pa, pdd pb, double r){
6.8 Half Plane Intersection
                                                                 if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
                                                                 if(abs(pb)<eps) return 0;</pre>
// NOTE: Point is complex<Real>
                                                                 double S, h, theta;
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
                                                                 double a=abs(pb),b=abs(pa),c=abs(pb-pa);
struct Line {
                                                                 double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
 Point st, ed;
                                                                 double cosC = dot(pa,pb) / a / b, C = acos(cosC);
 Point dir;
                                                                 if(a > r){
 Line (Point _s, Point _e)
                                                                 S = (C/2)*r*r;
   : st(_s), ed(_e), dir(_e - _s) {}
                                                                 h = a*b*sin(C)/c;
                                                                 if (h < r && B < PI/2)</pre>
                                                                   S = (acos(h/r)*r*r - h*sqrt(r*r-h*h));
bool operator<(const Line &lhs, const Line &rhs) {</pre>
 if (int cmp = argCmp(lhs.dir, rhs.dir))
    return cmp == -1;
                                                                 else if(b > r){
                                                                 theta = PI - B - asin(sin(B)/r*a);
  return ori(lhs.st, lhs.ed, rhs.st) < 0;</pre>
                                                                 S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
Point intersect(const Line &A, const Line &B) {
                                                                 else S = .5*sin(C)*a*b;
 Real t = cross(B.st - A.st, B.dir) /
   cross(A.dir, B.dir);
                                                                 return S;
  return A.st + t * A.dir;
                                                               double area_poly_circle(const vector<pdd> poly,
                                                                 const pdd &0,const double r){
                                                                 double S=0;
Real HPI(vector<Line> &lines) {
                                                                 for(int i=0;i<SZ(poly);++i)</pre>
  sort(lines.begin(), lines.end());
                                                                  S+=\_area(poly[i]-0,poly[(i+1)\%SZ(poly)]-0,r)
  deque<Line> que;
                                                                    *ori(0,poly[i],poly[(i+1)%SZ(poly)]);
  deque<Point> pt;
                                                                 return fabs(S);
  que.push_back(lines[0]);
  for (int i = 1; i < (int)lines.size(); i++) {
   if (argCmp(lines[i].dir, lines[i-1].dir) == 0)</pre>
                                                               6.12 Intersection of Two Circle
     continue;
                                                               bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
#define POP(L, R) \
    while (pt.size() > 0 \
                                                                pdd o1 = a.0, o2 = b.0;
                                                                 double r1 = a.R, r2 = b.R, d2 = abs2(o1 - o2),
      && ori(L.st, L.ed, pt.back()) < 0) \
      pt.pop_back(), que.pop_back(); \
                                                                     d = sqrt(d2);
    while (pt.size() > 0 \
                                                                if(d < max(r1, r2) - min(r1, r2) \mid \mid d > r1 + r2)
```

LL dx = x1-x2, dy = y1-y2;

```
return 0;
                                                                    return dx*dx+dy*dy;
pdd u = (o1 + o2) * 0.5
 + (o1 - o2) * ((r2 * r2 - r1 * r1) / (2 * d2));
                                                                   static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
double A = sqrt((r1 + r2 + d) * (r1 - r2 + d)
 * (r1 + r2 - d) * (-r1 + r2 + d));
                                                                   static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
                                                                   void init(vector<pair<int,int>> ip) {
pdd v = pdd(o1.Y - o2.Y, -o1.X + o2.X) * A
                                                                    n = ip.size();
  / (2 * d2);
                                                                    for (int i=0; i<n; i++) {</pre>
                                                                     tree[i].id = i;
p1 = u + v, p2 = u - v;
return 1;
                                                                     tree[i].x = ip[i].first;
                                                                     tree[i].y = ip[i].second;
6.13
      Tangent line of Two Circle
                                                                    root = build_tree(0, n-1, 0);
vector<Line> go(const Cir& c1,
                                                                   Node* build_tree(int L, int R, int d) {
 const Cir& c2, int sign1){
                                                                    if (L>R) return nullptr;
 // sign1 = 1 for outer tang, -1 for inter tang
                                                                    int M = (L+R)/2; tree[M].f = d%2;
vector<Line> ret;
                                                                    nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
double d_sq = norm2( c1.0 - c2.0 );
                                                                    tree[M].x1 = tree[M].x2 = tree[M].x;
if( d_sq < eps ) return ret;
double d = sqrt( d_sq );</pre>
                                                                    tree[M].y1 = tree[M].y2 = tree[M].y;
                                                                    tree[M].L = build_tree(L, M-1, d+1);
Pt v = (c2.0 - c1.0) / d;
                                                                    if (tree[M].L) {
double c = (c1.R - sign1 * c2.R) / d;
                                                                     tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
 if( c * c > 1 ) return ret;
double h = sqrt( max( 0.0 , 1.0 - c * c ) );
for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
                                                                     tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
 Pt n = \{ v.X * c - sign2 * h * v.Y ,
   v.Y * c + sign2 * h * v.X };
                                                                    tree[M].R = build_tree(M+1, R, d+1);
 Pt p1 = c1.0 + n * c1.R;
                                                                    if (tree[M].R) {
 Pt p2 = c2.0 + n * (c2.R * sign1);
                                                                     tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
 if( fabs( p1.X - p2.X ) < eps and
    fabs( p1.Y - p2.Y ) < eps )
                                                                     tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
   p2 = p1 + perp(c2.0 - c1.0);
                                                                     tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
 ret.push_back( { p1 , p2 } );
                                                                    return tree+M;
return ret;
                                                                   int touch(Node* r, int x, int y, LL d2){
                                                                    LL dis = sqrt(d2)+1;
6.14 Minimum Covering Circle
                                                                    if (x<r->x1-dis || x>r->x2+dis ||
template<typename P>
                                                                      y<r->y1-dis || y>r->y2+dis)
Circle getCircum(const P &a, const P &b, const P &c){
                                                                     return 0:
Real a1 = a.x-b.x, b1 = a.y-b.y;
                                                                    return 1;
Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
Real a2 = a.x-c.x, b2 = a.y-c.y;
                                                                   void nearest(Node* r,int x,int y,int &mID,LL &md2) {
Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
                                                                    if (!r || !touch(r, x, y, md2)) return;
Circle cc;
                                                                    LL d2 = dis2(r->x, r->y, x, y);
cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
                                                                    if (d2 < md2 | | (d2 == md2 && mID < r->id)) {
cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
                                                                     mID = r -> id;
cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
                                                                     md2 = d2:
return cc;
                                                                    // search order depends on split dim
                                                                    if ((r->f == 0 \&\& x < r->x) ||
template<typename P>
                                                                      (r->f == 1 \&\& y < r->y)) {
Circle MinCircleCover(const vector<P>& pts){
                                                                     nearest(r->L, x, y, mID, md2);
random_shuffle(pts.begin(), pts.end());
                                                                     nearest(r->R, x, y, mID, md2);
Circle c = { pts[0], 0 };
                                                                    } else {
for(int i=0;i<(int)pts.size();i++){</pre>
                                                                     nearest(r->R, x, y, mID, md2);
 if (dist(pts[i], c.o) <= c.r) continue;</pre>
                                                                     nearest(r->L, x, y, mID, md2);
 c = { pts[i], 0 };
for (int j = 0; j < i; j++) {
  if(dist(pts[j], c.o) <= c.r) continue;
  c.o = (pts[i] + pts[j]) / 2;</pre>
                                                                    }
                                                                   int query(int x, int y) {
                                                                    int id = 1029384756;
   c.r = dist(pts[i], c.o);
                                                                    LL d2 = 102938475612345678LL;
   for (int k = 0; k < j; k++) {
  if (dist(pts[k], c.o) <= c.r) continue;</pre>
                                                                    nearest(root, x, y, id, d2);
                                                                    return id:
    c = getCircum(pts[i], pts[j], pts[k]);
                                                                 } tree;
  }
                                                                        Rotating Sweep Line
return c;
                                                                 void rotatingSweepLine(pair<int, int> a[], int n) {
                                                                   vector<pair<int, int>> 1;
                                                                    1.reserve(n * (n - 1) / 2)
6.15 KDTree (Nearest Point)
                                                                    for (int i = 0; i < n; ++i)
const int MXN = 100005;
                                                                      for (int j = i + 1; j < n; ++j)
                                                                    l.emplace_back(i, j);
sort(l.begin(), l.end(), [&a](auto &u, auto &v){
  lld udx = a[u.first].first - a[u.second].first;
struct KDTree {
 struct Node {
  int x,y,x1,y1,x2,y2;
 int id,f;
Node *L, *R;
                                                                      lld udy = a[u.first].second - a[u.second].second;
lld vdx = a[v.first].first - a[v.second].first;
} tree[MXN], *root;
                                                                      11d vdy = a[v.first].second - a[v.second].second;
                                                                      if (udx == 0 or vdx == 0) return not udx == 0;
 int n;
LL dis2(int x1, int y1, int x2, int y2) {
                                                                      int s = sgn(udx * vdx);
```

return udy * vdx * s < vdy * udx * s;

```
vector<int> idx(n), p(n);
  iota(idx.begin(), idx.end(), 0);
sort(idx.begin(), idx.end(), [&a](int i, int j){
    return a[i] < a[j]; });
  for (int i = 0; i < n; ++i) p[idx[i]] = i;
  for (auto [i, j]: 1) {
    // do here
    swap(p[i], p[j]);
    idx[p[i]] = i, idx[p[j]] = j;
}
7
     Stringology
```

7.1 Hash

```
class Hash {
private:
  static constexpr int P = 127, Q = 1051762951;
  vector<int> h, p;
public:
 void init(const string &s){
  h.assign(s.size()+1, 0); p.resize(s.size()+1);
  for (size_t i = 0; i < s.size(); ++i)</pre>
   h[i + 1] = add(mul(h[i], P), s[i]);
  generate(p.begin(), p.end(), [x=1,y=1,this]()
     mutable{y=x;x=mul(x,P);return y;});
 int query(int 1, int r){ // 1-base (1, r]
   return sub(h[r], mul(h[1], p[r-1]));}
```

7.2 Suffix Array

```
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
int x[maxn], p[maxn], q[maxn * 2];
// sa[i]: sa[i]-th suffix is the \
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix \
// of suffix sa[i] and suffix sa[i - 1]
void pre(int *sa, int *c, int n, int z) {
 memset(sa, 0, sizeof(int) * n);
 memcpy(x, c, sizeof(int) * z);
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
  if (sa[i] && !t[sa[i] - 1])
   sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
 memcpy(x, c, sizeof(int) * z);
 for (int i = n - 1; i >= 0; --i)
  if (sa[i] && t[sa[i] - 1])
   sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q,
bool *t, int *c, int n, int z) {
 bool uniq = t[n - 1] = true;
 int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
 memset(c, 0, sizeof(int) * z);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
 for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
 if (uniq) {
  for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
  return;
 for (int i = n - 2; i \ge 0; --i)
  t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i)</pre>
  if (t[i] && !t[i - 1])
   sa[--x[s[i]]] = p[q[i] = nn++] = i;
 induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)
  if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
  bool neq = last < 0 || \</pre>
   memcmp(s + sa[i], s + last,
(p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
  ns[q[last = sa[i]]] = nmxz += neq;
 }}
```

```
sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
 pre(sa, c, n, z);
for (int i = nn - 1; i >= 0; --i)
  sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
 induce(sa, c, s, t, n, z);
void build(const string &s) {
 for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
 _s[(int)s.size()] = 0; // s shouldn't contain 0
 sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];</pre>
 for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
 int ind = 0; hi[0] = 0;
 for (int i = 0; i < (int)s.size(); ++i) {</pre>
  if (!rev[i]) {
   ind = 0:
   continue;
  while (i + ind < (int)s.size() && \</pre>
   s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  hi[rev[i]] = ind ? ind-- : 0;
}}
7.3
     Suffix Automaton
```

```
struct Node{
Node *green, *edge[26];
 int max_len;
 Node(const int _max_len)
 : green(NULL), max_len(_max_len){
  memset(edge,0,sizeof(edge));
} *ROOT, *LAST;
void Extend(const int c) {
Node *cursor = LAST:
 LAST = new Node((LAST->max_len) + 1);
 for(;cursor&!cursor->edge[c]; cursor=cursor->green)
 cursor->edge[c] = LAST;
 if (!cursor)
 LAST->green = ROOT;
 else {
  Node *potential_green = cursor->edge[c];
  if((potential_green->max_len)==(cursor->max_len+1))
   LAST->green = potential_green;
  else {
//assert(potential_green->max_len>(cursor->max_len+1));
   Node *wish = new Node((cursor->max_len) + 1);
   for(;cursor && cursor->edge[c]==potential_green;
      cursor = cursor->green)
    cursor->edge[c] = wish;
   for (int i = 0; i < 26; i++)
    wish->edge[i] = potential_green->edge[i];
   wish->green = potential_green->green;
   potential_green->green = wish;
   LAST->green = wish;
char S[10000001], A[10000001];
int N;
int main(){
scanf("%d%s", &N, S);
ROOT = LAST = new Node(0);
 for (int i = 0; S[i]; i++)
 Extend(S[i] - 'a');
 while (N--){
  scanf("%s", A);
  Node *cursor = ROOT;
  bool ans = true;
  for (int i = 0; A[i]; i++){
   cursor = cursor->edge[A[i] - 'a'];
   if (!cursor) {
   ans = false;
    break;
 puts(ans ? "Yes" : "No");
 return 0;
```

7.4 KMP

```
vector<int> kmp(const string &s) {
 vector<int> f(s.size(), 0);
 /* f[i] = length of the longest prefix
   (excluding s[0:i]) such that it coincides with the suffix of s[0:i] of the same length */
 /* i + 1 - f[i] is the length of the
   smallest recurring period of s[0:i] */
 int k = 0:
 for (int i = 1; i < (int)s.size(); ++i) {</pre>
  while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
  if (s[i] == s[k]) ++k;
  f[i] = k;
 return f;
vector<int> search(const string &s, const string &t) {
 // return 0-indexed occurrence of t in s
 vector<int> f = kmp(t), r;
 for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
 while(k > 0 && (k==(int)t.size() \mid \mid s[i]!=t[k]))
   k = f[k - 1];
  if (s[i] == t[k]) ++k;
  if (k == (int)t.size()) r.push_back(i-t.size()+1);
return res;
7.5 Z value
char s[MAXN];
int len,z[MAXN];
void Z_value() {
 int i, j, left, right;
 z[left=right=0]=len;
 for(i=1;i<len;i++)</pre>
  j=max(min(z[i-left], right-i),0);
  for(;i+j<len&&s[i+j]==s[j];j++)</pre>
  if(i+(z[i]=j)>right)right=i+z[left=i];
7.6 Manacher
int z[maxn];
int manacher(const string& s) {
  string t = ".";
 for(char c: s) t += c, t += '.';
 int 1 = 0, r = 0, ans = 0;
 for (int i = 1; i < t.length(); ++i) {</pre>
 z[i] = (r > i'] min(z[2 * 1' - i], r - i) : 1);
while (i - z[i] >= 0 && i + z[i] < t.length()) {
   if(t[i - z[i]] == t[i + z[i]]) ++z[i];
   else break;
  if (i + z[i] > r) r = i + z[i], l = i;
 for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);</pre>
 return ans;
7.7 Lexico Smallest Rotation
string mcp(string s){
 int n = s.length();
 s += s;
 int i=0, j=1;
 while (i<n && j<n){</pre>
  int k = 0;
  while (k < n \&\& s[i+k] == s[j+k]) k++;
  if (s[i+k] <= s[j+k]) j += k+1;</pre>
 else i += k+1;
  if (i == j) j++;
 int ans = i < n ? i : j;</pre>
 return s.substr(ans, n);
```

7.8 BWT

#define SIGMA 26

#define BASE 'a'

struct BurrowsWheeler{

vector<int> v[SIGMA];

void BWT(char* ori, char* res){
 // make ori -> ori + ori

```
void iBWT(char* ori, char* res){
  for( int i = 0 ; i < SIGMA ; i ++ )</pre>
   v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )</pre>
   v[ ori[i] - BASE ].push_back( i );
  vector<int> a:
  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )
for( auto j : v[ i ] ){</pre>
    a.push_back( j );
    ori[ ptr ++ ] = BASE + i;
  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
   ptr = a[ ptr ];
  res[ len ] = 0;
} bwt;
7.9 Palindromic Tree
struct palindromic_tree{
 struct node{
  int next[26],f,len;
  int cnt,num,st,ed;
  node(int l=0):f(0),len(1),cnt(0),num(0) {
   memset(next, 0, sizeof(next)); }
 vector<node> st:
 vector<char> s;
 int last,n;
 void init(){
  st.clear();s.clear();last=1; n=0;
  st.push_back(0);st.push_back(-1);
  st[0].f=1;s.push_back(-1); }
 int getFail(int x){
  while(s[n-st[x].len-1]!=s[n])x=st[x].f;
  return x;}
 void add(int c){
  s.push_back(c-='a'); ++n;
  int cur=getFail(last);
  if(!st[cur].next[c]){
   int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
   st[now].num=st[st[now].f].num+1;
  last=st[cur].next[c];
  ++st[last].cnt;}
 void dpcnt() {
  for (int i=st.size()-1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
 int size(){ return st.size()-2;}
} pt;
int main() {
 string s; cin >> s; pt.init();
 for (int i=0; i<SZ(s); i++) {</pre>
  int prvsz = pt.size(); pt.add(s[i]);
  if (prvsz != pt.size()) {
   int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [l,r]: s.substr(l, r-l+1)
  }
 return 0;
8
     Misc
```

// then build suffix array

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\ldots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + d_2 + \ldots + d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all 1 < k < n.

8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W\subseteq X, |W|\leq |N_G(W)|$

8.1.7 Euler's planar graph formula

$$V - E + F = C + 1, E \le 3V - 6$$
(?)

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Lucas's theorem

 ${m\choose n}\equiv\prod_{i=0}^k{m_i\choose n_i}\pmod{p}, \text{ where } m=m_kp^k+m_{k-1}p^{k-1}+\cdots+m_1p+m_0,$ and $n=n_kp^k+n_{k-1}p^{k-1}+\cdots+n_1p+n_0.$

8.1.10 Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2)$, find maximum $S\in I_1\cap I_2$. For each iteration, build the directed graph and find a shortest path from s to t.

- $s \rightarrow x : S \sqcup \{x\} \in I_1$
- $x \to t : S \sqcup \{x\} \in I_2$
- $y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1$ (y is in the unique circuit of $S \sqcup \{x\}$)
- $x \to y: S \setminus \{y\} \sqcup \{x\} \in I_2$ (y is in the unique circuit of $S \sqcup \{x\}$)

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|$. In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

8.2 DP-opt Condition

8.2.1 totally monotone (concave/convex)

 $\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}$

8.2.2 monge condition (concave/convex)

 $\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}$

8.3 Convex 1D/1D DP

```
struct segment {
  int i, 1, r;
  segment() {}
  segment(int a, int b, int c): i(a), l(b), r(c) {}
};
  inline lld f(int l, int r){return dp[l] + w(l+1, r);}
  void solve() {
   dp[0] = 0;
   deque<segment> dq; dq.push_back(segment(0, 1, n));
   for (int i = 1; i <= n; ++i) {
    dp[i] = f(dq.front().i, i);
}</pre>
```

```
while(dq.size()&&dq.front().r<i+1) dq.pop_front();
dq.front().l = i + 1;
segment seg = segment(i, i + 1, n);
while (dq.size() &&
    f(i, dq.back().l)<f(dq.back().i, dq.back().l))
        dq.pop_back();
if (dq.size()) {
    int d = 1 << 20, c = dq.back().l;
    while (d >>= 1) if (c + d <= dq.back().r)
        if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
        dq.back().r = c; seg.l = c + 1;
    }
    if (seg.l <= n) dq.push_back(seg);
}</pre>
```

8.4 ConvexHull Optimization

```
struct Line {
mutable int64_t a, b, p;
 bool operator<(const Line &rhs) const { return a < rhs
    .a; }
bool operator<(int64_t x) const { return p < x; }</pre>
};
struct DynamicHull : multiset<Line, less<>> {
 static const int64_t kInf = 1e18;
 bool Isect(iterator x, iterator y)
  auto Div = [](int64_t a, int64_t b) {
    return a / b - ((a ^ b) < 0 && a % b); }
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(int64_t a, int64_t b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
  if (x != begin() \&\& Isect(--x, y)) Isect(x, y = erase)
    (y));
  while ((y = x) != begin() && (--x)->p >= y->p) Isect(
    x, erase(y));
 int64_t Query(int64_t x) {
  auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
 }
};
```

8.5 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

8.6 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u)
 dfn[u]=low[u]=++dfs_idx;
 for(int i=0;i<(int)init_g[u].size();i++){</pre>
  int v=init_g[u][i];
  if(v==par[u]) continue;
  if(!dfn[v]){
   par[v]=u;
   tarjan(v);
   low[u]=min(low[u],low[v]);
   if(dfn[u]<low[v]){</pre>
    g[u].push_back(v);
    g[v].push_back(u);
  }else{
   low[u]=min(low[u],dfn[v]);
   if(dfn[v]<dfn[u]){</pre>
```

```
int temp_v=u;
                                                               const static int maxnode=210*210;
    bcc_id++;
                                                               int n, m, size, row[maxnode], col[maxnode];
                                                               int U[maxnode], D[maxnode], L[maxnode], R[maxnode];
    while(temp_v!=v){
     g[bcc_id+n].push_back(temp_v);
                                                               int H[maxn], S[maxm], ansd, ans[maxn];
     g[temp_v].push_back(bcc_id+n);
                                                               void init(int _n, int _m) {
     temp_v=par[temp_v];
                                                                n = _n, m = _m;
                                                                for(int i = 0; i <= m; ++i) {</pre>
                                                                S[i] = 0;
    g[bcc_id+n].push_back(v);
    g[v].push_back(bcc_id+n);
                                                                 U[i] = D[i] = i;
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
                                                                L[i] = i-1, R[i] = i+1;
                                                                R[L[0] = size = m] = 0;
                                                                for(int i = 1; i <= n; ++i) H[i] = -1;
 }
int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
                                                               void Link(int r, int c) {
void dfs(int u,int fa){
                                                                ++S[col[++size] = c];
                                                                row[size] = r; D[size] = D[c];
 if(u<=n){
                                                                U[D[c]] = size; U[size] = c; D[c] = size;
  for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                if(H[r] < 0) H[r] = L[size] = R[size] = size;
   int v=g[u][i];
   if(v==fa) continue;
                                                                else {
                                                                R[size] = R[H[r]];
   dfs(v,u);
   memset(tp,0x8f,sizeof tp);
                                                                 L[R[H[r]]] = size;
                                                                 L[size] = H[r];
   if(v<=n){
    tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
                                                                R[H[r]] = size;
    tp[1]=max(
                                                               }
     dp[u][0]+dp[v][0]+1
     dp[u][1]+max(dp[v][0],dp[v][1])
                                                               void remove(int c) {
                                                                L[R[c]] = L[c]; R[L[c]] = R[c];
                                                               for(int i = D[c]; i != c; i = D[i])
for(int j = R[i]; j != i; j = R[j]) {
   }else{
    tp[0]=dp[u][0]+dp[v][0];
                                                                  U[D[j]] = U[j];
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
                                                                  D[U[j]] = D[j];
   dp[u][0]=tp[0],dp[u][1]=tp[1];
                                                                  --S[col[j]];
                                                                 }
 }else{
  for(int i=0;i<(int)g[u].size();i++){</pre>
                                                               void resume(int c) {
   int v=g[u][i];
                                                               L[R[c]] = c; R[L[c]] = c;
                                                                for(int i = U[c]; i != c; i = U[i])
   if(v==fa) continue;
   dfs(v,u);
                                                                 for(int j = L[i]; j != i; j = L[j]) {
                                                                 U[D[j]] = j;
  min_dp[0][0]=0;
                                                                  D[U[j]] = j
  min_dp[1][1]=1;
                                                                  ++S[col[j]];
  min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
                                                                 }
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
                                                               void dance(int d) {
                                                                if(d>=ansd) return;
   if(v==fa) continue;
   memset(tmp,0x8f,sizeof tmp);
                                                                if(R[0] == 0) {
                                                                 ansd = d;
   tmp[0][0]=max(
    min_dp[0][0]+max(dp[v][0],dp[v][1]),
                                                                 return;
    min_dp[0][1]+dp[v][0]
                                                                int c = R[0];
                                                                for(int i = R[0]; i; i = R[i])
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
   tmp[1][0]=max(
                                                                if(S[i] < S[c]) c = i;
    \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
                                                                remove(c);
    min_dp[1][1]+dp[v][0]
                                                                for(int i = D[c]; i != c; i = D[i]) {
                                                                 ans[d] = row[i]
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
                                                                 for(int j = R[i]; j != i; j = R[j])
   memcpy(min_dp,tmp,sizeof tmp);
                                                                  remove(col[j]);
                                                                 dance(d+1);
  dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
                                                                 for(int j = L[i]; j != i; j = L[j])
  dp[u][0]=min_dp[0][0];
                                                                  resume(col[j]);
                                                                resume(c);
int main(){
                                                               }
 int m,a,b;
                                                             } sol;
 scanf("%d%d",&n,&m);
                                                                   Tree Knapsack
 for(int i=0;i<m;i++){
  scanf("%d%d",&a,&b);</pre>
                                                             int dp[N][K];PII obj[N];
  init_g[a].push_back(b);
                                                             vector<int> G[N];
  init_g[b].push_back(a);
                                                             void dfs(int u, int mx){
                                                              for(int s: G[u]) {
                                                                if(mx < obj[s].first) continue;</pre>
 par[1]=-1;
 tarjan(1);
                                                                for(int i=0;i<=mx-obj[s].FF;i++)</pre>
                                                                 dp[s][i] = dp[u][i];
 dfs(1,-1);
 printf("%d\n", max(dp[1][0], dp[1][1]));
                                                                dfs(s, mx - obj[s].first);
 return 0;
                                                                for(int i=obj[s].FF;i<=mx;i++)</pre>
                                                                 dp[u][i] = max(dp[u][i],
                                                                  dp[s][i - obj[s].FF] + obj[s].SS);
8.7 DLX
struct DLX {
 const static int maxn=210;
                                                             int main(){
                                                              int n, k; cin >> n >> k;
const static int maxm=210;
```

```
for(int i=1;i<=n;i++){</pre>
  int p; cin >> p;
  G[p].push_back(i);
  cin >> obj[i].FF >> obj[i].SS;
 dfs(0, k); int ans = 0;
 for(int i=0;i<=k;i++) ans = max(ans, dp[0][i]);
 cout << ans << '\n';
 return 0;
8.9 N Queens Problem
vector< int > solve( int n ) {
 // no solution when n=2, 3
 vector< int > ret;
if ( n % 6 == 2 ) {
  for ( int i = 2 ; i <= n ; i += 2 )
    ret.push_back( i );</pre>
  ret.push_back( 3 ); ret.push_back( 1 );
for ( int i = 7 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 5 );
 } else if ( n % 6 == 3 ) {
  for ( int i = 4 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
  ret.push_back( 2 );
  for ( int i = 5 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
  ret.push_back( 1 ); ret.push_back( 3 );
 } else {
for ( int i = 2 ; i <= n ; i += 2 )</pre>
   ret.push_back( i );
  for ( int i = 1 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
 return ret;
8.10 Aliens Optimization
long long Alien() {
 long long c = kInf;
 for (int d = 60; d >= 0; --d) {
  // cost can be negative, depending on the problem.
  if (c - (1LL << d) < 0) continue;</pre>
  long long ck = c - (1LL \ll d);
  pair<long long, int> r = check(ck);
  if (r.second == k) return r.first - ck * k;
  if (r.second < k) c = ck;</pre>
pair<long long, int> r = check(c);
return r.first - c * k;
8.11
      To Check When Submit
   · Array out of bound.
   • long long / double cast.
```

- Initialization.
- Delete all debug code.
- Check sample testcase after simple modification.