Contents

```
1 Basic
                               5.26 Adaptive Simpson . . 15
1.1
     vimrc .....
                               5.27 Golden Ratio Search.
                                                        15
     Debug Macro.....
1.2
                              6 Geometry
                                                        15
1.3
     Increase Stack .....
                                   Basic Geometry....
                                                        15
                               6.1
     Pragma Optimization
1.4
                                   2D Convex Hull .....
                                                        16
1.5
     IO Optimization....
                               6.3
                                   2D Farthest Pair ....
                                                        16
     SVG Writer .....
1.6
                                   MinMax Enclosing
2 Data Structure
                                   16
2.1
     Dark Magic . . . . . . . . . .
                          2
                                   Minkowski Sum..... 16
                                   Segment Intersection 16
     Link-Cut Tree . . . . . .
2.2
                                   Half Plane Intersec-
2.3
     LiChao Segment Tree
                               6.7
                                                        16
                                   tion .....
2.4 Treap .....
                                   SegmentDist
                               6.8
2.5 Linear Basis ......
                          3
                                                        16
                                   (Sausage) . . . . . . . . . .
     Binary Search On
2.6
                                   Rotating Sweep Line
                                                        16
                               6.9
     3
                               6.10 Polygon Cut . . . . . . . .
                          3
3 Graph
                                   Point In Simple
3.1
     SCC .....
                                   Polygon............ 17
3.2
3.3
     2-SAT .....
                               6.12 Point In Hull (Fast) .. 17
     BCC ..
                               6.13 Point In Polygon
3.4
     Round Square Tree..
                                                       17
3.5
     Edge TCC .....
                                   6.14 Cyclic Ternary Search 17
3.6
     DMST.....
     Dominator Tree . . . . .
                               6.15 Tangent of Points
3.7
3.8
    Edge Coloring .....
                                   To Hull . . . . . . . . . . . . . . . .
                               6.16 Circle Class & Inter-
     Centroid Decomposition.....
3.9
                                   section .....
                               6.17 Circle Common Tangent .....
3.10 Lowbit Decompo-
     sition .........
                               6.18 Line-Circle Inter-
3.11
     Virtual Tree.....
                                   section ......
3.12 Tree Hashing ......
                               6.19 Poly-Circle Inter-
3.13 Mo's Algorithm on
                                   section .....
6.20 Minimum Covering
                                   Circle .....
3.15 MaximalClique .....
                               6.21 Circle Union . . . . . . .
                                                        18
                               6.22 Polygon Union..... 18
3.16 MaximumClique ....
                               6.23 3D Point .....
                                                        19
    Minimum Mean
Cycle .....
                               6.24 3D Convex Hull .....
                                                        19
                               6.25 3D Projection . . . . . .
                                                        19
4 Flow & Matching
                          7
                              6.26 3D Skew Line Near-
est Point .....
    HopcroftKarp.....
                                                        19
     Kuhn Munkres.....
                               6.27 Delaunay .....
                                                        19
4.3
     Flow Models.....
                               6.28 Build Voronoi . . . . . .
                                                       20
                               6.29 kd Tree (Nearest
     HLPP .....
 4.5
                                   Point) .....
 4.6
     Global Min-Cut .....
                               6.30 kd Closest Pair (3D
 4.7
     GomoryHu Tree . . . . .
                                   ver.) .....
     Minimum Cost Max Flow.....
4.8
                               6.31 Simulated Annealing
                                                       20
                          9
4.9 Dijkstra Cost Flow . . .
                               6.32 Triangle Centers ....
                                                        21
4.10 Minimum Cost Cir-
                                                        21
                              7 Stringology
     culation . . . . . . . .
                                   Hash .....
     General
               Graph
                         10
                               7.2
                                   Suffix Array .....
     Matching .....
4.12 Weighted Matching . 10
                                   Suffix Array Tools . . . .
                                   Ex SAM .....
5 Math
                                   KMP.....
5.1
     Common Bounds ...
                          11
                               7.6
7.7
                                   Z value .....
5.2
     Equations . . . . . . . . . . . .
                                   Manacher . . . . . . . . . . .
                                                        22
5.3
     Extended FloorSum.
                          11
                               7.8
                                   Lyndon Factorization 22
5.4
     Integer Division . . . . .
                               7.9
                                   Main Lorentz .....
                               7.10 BWT.....
                                                        22
5.5
     FloorSum .....
                               7.11 Palindromic Tree....
                                                       22
5.6
     ModMin .....
     Floor Monoid Product
                                                       23
                              8 Misc
     ax+by=gcd .....
                                   Theorems .....
                               8.1
5.9 Chinese Remainder .
                                   Weight Matroid In-
5.10 DiscreteLog ......
                                   tersection .....
                                                       23
5.11 Quadratic Residue ...
                         12
                                   Stable Marriage . . . . .
                                                       23
5.12 FWT.....
                                   Bitset LCS .....
                                                       24
5.13 Packed FFT .....
                               8.5 Prefix Substring LCS. 24
5.14 CRT for arbitrary mod 12
                               8.6 Convex 1D/1D DP .... 24
5.15 NTT/FFT .....
                         13
                                   ConvexHull Opti-
5.16 Formal Power Series
                                   mization .....
5.17 Partition Number ...
                                   Min Plus Convolution 24
5.18 Pi Count .....
                         14
                               8.9 De-Bruijn .....
                                                       24
5.19 Miller Rabin .....
                               8.10 Josephus Problem ...
5.20 Pollard Rho ......
                               8.11 N Queens Problem . .
5.21 Berlekamp Massey . .
                                                       24
                               8.12 Tree Knapsack ..... 24
5.22 Gauss Elimination...
                               8.13 Manhattan MST..... 24
5.23 Charateristic Poly-
                               8.14 Binary Search On
     nomial . . . . . . . . . . . . . . . . . .
                                   Fraction....
5.24 Simplex . . . . . . . . . . . .
                         15
                                   Barrett Reduction . . .
5.25 Simplex Construction 15
                               8.16 Montgomery .....
```

```
Basic
1.1
     vimrc
se is nu ru et tgc sc hls cin cino+=j1 sw=2 sts=2 bs=2
    mouse=a "encoding=utf-8 ls=2
syn on | colo desert | filetype indent on
inoremap {<CR> {<CR>}<ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -g -std=gnu++20 -</pre>
    DCKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
    Wconversion -fsanitize=address,undefined,float-
    divide-by-zero, float-cast-overflow && echo success<
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -g -std=gnu</pre>
    ++20 && echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space
    :]' \| md5sum \| cut -c-6
let c_no_curly_error=1
" setxkbmap -option caps:ctrl_modifier
1.2 Debug Macro [a45c59]
#define all(x) begin(x), end(x)
#ifdef CKISEKI
#include <experimental/iterator>
#define safe cerr<<__PRETTY_FUNCTION__<<" line "<</pre>
     __LINE__<<" safe\n"
#define debug(a...) debug_(#a, a)
#define orange(a...) orange_(#a, a)
void debug_(auto s, auto ...a) {
  cerr << "\e[1;32m(" << s << ") = (";</pre>
 int f = 0;
 (..., (cerr << (f++ ? ", " : "") << a));
 cerr << ")\e[0m\n";</pre>
void orange_(auto s, auto L, auto R) {
  cerr << "\e[1;33m[ " << s << " ] = [ ";</pre>
 using namespace experimental;
 copy(L, R, make_ostream_joiner(cerr, ", "));
cerr << " ]\e[0m\n";</pre>
#else
#define safe ((void)0)
#define debug(...) safe
#define orange(...) safe
#endif
1.3 Increase Stack [b6856c]
const int size = 256 << 20;</pre>
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
       _("movq %0, %%rsp\n"::"r"(bak));
  asm
1.4 Pragma Optimization [6006f6]
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popent,abm,mmx,avx,arch=skylake")
  builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
1.5 IO Optimization [c9494b]
static inline int gc() {
 constexpr int B = 1<<20; static char buf[B], *p, *q;</pre>
 if (p == q) q = (p = buf) + fread(buf, 1, B, stdin);
 return q == buf ? EOF : *p++;
1.6 SVG Writer [57436c]
class SVG {
 void p(string_view s) { o << s; }</pre>
 void p(string_view s, auto v, auto... vs) {
  auto i = s.find('$');
  o << s.substr(0, i) << v, p(s.substr(i + 1), vs...);
 ofstream o; string c = "red";
public:
 SVG(auto f,auto x1,auto y1,auto x2,auto y2) : o(f) {
  p("<svg xmlns='http://www.w3.org/2000/svg'</pre>
   "viewBox='$ $ $'>\n"
   "<style>*{stroke-width:0.5%;}</style>\n",
 x1, -y2, x2 - x1, y2 - y1); }
~SVG() { p("</svg>\n"); }
 SVG &color(string nc) { return c = nc, *this; }
 void line(auto x1, auto y1, auto x2, auto y2) {
 p("<line x1='$' y1='$' x2='$' y2='$' stroke='$'/>\n",
   x1, -y1, x2, -y2, c); }
```

```
void circle(auto x, auto y, auto r) {
p("<circle cx='$' cy='$' r='$' stroke='$' "</pre>
                                                             void access(int x) {
   "fill='none'/>\n", x, -y, r, c); }
                                                              for (int u = x, last = 0; u; u = cur.pa) {
void text(auto x, auto y, string s, int w = 12) {
                                                               splay(u);
 p("<text x='$' y='$' font-size='$px'>$</text>\n",
                                                               cur.vir = cur.vir + o[rc].sub - o[last].sub;
  x, -y, w, s); }
                                                               rc = last; up(last = u);
                                                              splay(x);
2
     Data Structure
2.1 Dark Magic [095f25]
                                                             int find_root(int u) {
#include <ext/pb_ds/assoc_container.hpp>
                                                              int la = 0;
#include <ext/pb_ds/priority_queue.hpp>
                                                              for (access(u); u; u = lc) down(la = u);
using namespace __gnu_pbds;
                                                              return la;
// heap tags: paring/binary/binomial/rc_binomial/thin
template<typename T>
                                                             void split(int x, int y) { chroot(x); access(y); }
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>, \
                                                             void chroot(int u) { access(u); set_rev(u); }
                  pairing_heap_tag>;
                                                            public:
// pbds_heap::point_iterator
                                                             LCT(int n = 0) : o(n + 1) {}
  x = pq.push(10); pq.modify(x, 87); a.join(b);
                                                             int add(const Val &v = {}) {
// tree tags: rb_tree_tag/ov_tree_tag/splay_tree_tag
                                                              o.push_back(v);
template<typename T>
                                                              return int(o.size()) - 2;
using ordered_set = tree<T, null_type, less<T>,
   rb_tree_tag, tree_order_statistics_node_update>;
                                                             void set_val(int u, const Val &v) {
// find_by_order, order_of_key
                                                              splay(++u); cur.v = v; up(u);
// hash tables: cc_hash_table/gp_hash_table
2.2 Link-Cut Tree [60627f]
                                                             void set_sval(int u, const SVal &v) {
template <typename Val, typename SVal> class LCT {
                                                             access(++u); cur.sv = v; up(u);
struct node {
                                                             Val query(int x, int y) {
 int pa, ch[2];
 bool rev;
                                                              split(++x, ++y); return o[y].prod;
 Val v, prod, rprod;
 SVal sv, sub, vir;
                                                             SVal subtree(int p, int u) {
 node() : pa{0}, ch{0, 0}, rev{false}, v{}, prod{},
                                                              chroot(++p); access(++u);
                                                              return cur.vir + cur.sv;
    rprod{}, sv{}, sub{}, vir{} {};
#define cur o[u]
                                                             bool connected(int u, int v) {
#define lc cur.ch[0]
                                                              return find_root(++u) == find_root(++v); }
#define rc cur.ch[1]
                                                             void link(int x, int y) {
                                                              chroot(++x); access(++y);
vector<node> o:
                                                              o[y].vir = o[y].vir + o[x].sub;
bool is_root(int u) const {
 return o[cur.pa].ch[0] != u && o[cur.pa].ch[1] != u;
                                                              up(o[x].pa = y);
                                                             void cut(int x, int y) {
bool is_rch(int u) const {
                                                              split(++x, ++y);
 return o[cur.pa].ch[1] == u && !is_root(u); }
void down(int u) {
                                                              o[y].ch[0] = o[x].pa = 0; up(y);
 if (not cur.rev) return;
                                                            #undef cur
 if (lc) set_rev(lc);
                                                            #undef lc
 if (rc) set_rev(rc);
                                                            #undef rc
 cur.rev = false;
                                                           };
}
void up(int u) {
                                                            2.3 LiChao Segment Tree [b9c827]
 cur.prod = o[lc].prod * cur.v * o[rc].prod;
 cur.rprod = o[rc].rprod * cur.v * o[lc].rprod;
                                                            struct L {
 cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
                                                             int m, k, id;
                                                             L(): id(-1) {}
void set_rev(int u) {
                                                             L(int a, int b, int c) : m(a), k(b), id(c) {}
 swap(lc, rc), swap(cur.prod, cur.rprod);
                                                             int at(int x) { return m * x + k; }
 cur.rev ^= 1;
                                                            class LiChao {
void rotate(int u) {
                                                           private:
 int f = cur.pa, g = o[f].pa, l = is_rch(u);
if (cur.ch[l ^ 1]) o[cur.ch[l ^ 1]].pa = f;
                                                             int n; vector<L> nodes;
                                                             static int lc(int x) { return 2 * x + 1; }
 if (not is_root(f)) o[g].ch[is_rch(f)] = u;
                                                             static int rc(int x) { return 2 * x + 2; }
 o[f].ch[l] = cur.ch[l ^ 1];
                                                             void insert(int l, int r, int id, L ln) {
 cur.ch[l ^ 1] = f;
                                                              int m = (l + r) >> 1;
 cur.pa = g, o[f].pa = u;
                                                              if (nodes[id].id == -1)
 up(f);
                                                               return nodes[id] = ln, void();
                                                              bool atLeft = nodes[id].at(l) < ln.at(l);</pre>
                                                              if (nodes[id].at(m) < ln.at(m))</pre>
void splay(int u) {
 vector<int> stk = {u};
                                                               atLeft ^= 1, swap(nodes[id], ln);
                                                              if (r - l == 1) return;
 while (not is_root(stk.back()))
                                                              if (atLeft) insert(l, m, lc(id), ln);
   stk.push_back(o[stk.back()].pa);
 while (not stk.empty()) {
                                                              else insert(m, r, rc(id), ln);
  down(stk.back());
  stk.pop_back();
                                                             int query(int l, int r, int id, int x) {
                                                              int m = (l + r) >> 1, ret = 0;
                                                              if (nodes[id].id != -1) ret = nodes[id].at(x);
 for (int f = cur.pa; not is_root(u); f = cur.pa) {
  if (!is_root(f))
                                                              if (r - l == 1) return ret;
   rotate(is_rch(u) == is_rch(f) ? f : u);
                                                              if (x < m) return max(ret, query(l, m, lc(id), x));</pre>
  rotate(u);
                                                              return max(ret, query(m, r, rc(id), x));
 up(u);
                                                           public:
```

```
LiChao(int n_{-}): n(n_{-}), nodes(n * 4) {}
 void insert(L ln) { insert(0, n, 0, ln); }
                                                                 return l + 1 - sz;
 int query(int x) { return query(0, n, 0, x); }
                                                                } else sum = s, l++;
                                                               } while (lowbit(l) != l);
};
                                                               return n + 1:
2.4
      Treap [ae576c]
__gnu_cxx::sfmt19937 rnd(7122); // <ext/random>
                                                              int find_last(int r, auto &&check) {
namespace Treap {
                                                               if (r <= 0) return -1;
struct node {
                                                               r += sz; push(r - 1); Monoid sum; // identity
int size, pri; node *lc, *rc, *pa;
node() : size(1), pri(rnd()), lc(0), rc(0), pa(0) {}
                                                               do {
 void pull() {
                                                                while (r > 1 and (r & 1)) r >>= 1;
 size = 1; pa = 0;
                                                                if (auto s = nd[r] + sum; check(s)) {
  if (lc) { size += lc->size; lc->pa = this; }
                                                                 while (r < sz) {</pre>
 if (rc) { size += rc->size; rc->pa = this; }
                                                                  prop(r); r = (r << 1) | 1;
 }
                                                                   if (auto nxt = nd[r] + sum; not check(nxt))
                                                                    sum = nxt, r--;
int SZ(node *x) { return x ? x->size : 0; }
node *merge(node *L, node *R) {
  if (not L or not R) return L ? L : R;
                                                                 return r - sz;
                                                                } else sum = s;
 if (L->pri > R->pri)
                                                               } while (lowbit(r) != r);
  return L->rc = merge(L->rc, R), L->pull(), L;
                                                               return -1;
  return R->lc = merge(L, R->lc), R->pull(), R;
                                                                   Graph
                                                                   SCC [1fdaba]
void splitBySize(node *o, int k, node *&L, node *&R) {
                                                              3.1
 if (not o) L = R = 0;
                                                              class SCC { // test @ library checker
 else if (int s = SZ(o->lc) + 1; s <= k)
                                                              protected:
  L=o, splitBySize(o->rc, k-s, L->rc, R), L->pull();
                                                               int n, dfc, nscc; vector<vector<int>> G;
 else
                                                               vector<int> vis, low, idx, stk;
  R=o, splitBySize(o->lc, k, L, R->lc), R->pull();
                                                               void dfs(int i) {
\frac{1}{N} = \frac{1}{N} \sum_{i=1}^{N} \frac{1}{N} \sum_{i=1}^{N} \frac{1}{N} 
                                                                vis[i] = low[i] = ++dfc;
int getRank(node *o) { // 1-base
                                                                stk.push_back(i);
 int r = SZ(o->lc) + 1;
                                                                for (int j : G[i]) {
                                                                 if (!vis[j]) {
 for (; o->pa; o = o->pa)
  if (o->pa->rc == o) r += SZ(o->pa->lc) + 1;
                                                                  dfs(j);
 return r;
                                                                  low[i] = min(low[i], low[j]);
                                                                 } else if (vis[j] != -1) {
} // namespace Treap
                                                                   low[i] = min(low[i], vis[j]);
2.5 Linear Basis [138d5d]
template <int BITS, typename S = int> struct Basis {
                                                                if (low[i] == vis[i]) {
 static constexpr S MIN = numeric_limits<S>::min();
                                                                 while (true) {
 array<pair<llu, S>, BITS> b;
                                                                  int x = stk.back(); stk.pop_back();
 Basis() { b.fill({0, MIN}); }
                                                                   idx[x] = nscc; vis[x] = -1;
 void add(llu x, S p) {
                                                                   if (x == i) break;
  for (int i = BITS-1; i>=0; i--) if (x >> i & 1) {
                                                                 }
   if (b[i].first == 0) return b[i]={x, p}, void();
   if (b[i].second < p)</pre>
                                                                 ++nscc;
                                                                }
    swap(b[i].first, x), swap(b[i].second, p);
                                                               7
   x ^= b[i].first;
                                                              public:
  }
                                                               SCC(int n_{-}) : n(n_{-}), dfc(0), nscc(0), G(n),
 optional<llu> query_kth(llu v, llu k) {
                                                                vis(n), low(n), idx(n) {}
                                                               void add_edge(int u, int v) { G[u].push_back(v); }
 vector<pair<llu, int>> o;
                                                               void solve() {
  for (int i = 0; i < BITS; i++)</pre>
                                                                for (int i = 0; i < n; i++) if (!vis[i]) dfs(i); }</pre>
   if (b[i].first) o.emplace_back(b[i].first, i);
                                                               int get_id(int x) { return idx[x]; }
  if (k >= (1ULL << o.size())) return {};</pre>
                                                               int count() { return nscc; }
  for (int i = int(o.size()) - 1; i >= 0; i--)
   if ((k >> i & 1) ^ (v >> o[i].second & 1))
                                                               3.2
                                                                   2-SAT [ca961f]
    v ^= o[i].first;
                                                              struct TwoSat : SCC {
  return v:
                                                               void orr(int x, int y) {
 Basis filter(S l) {
                                                                if ((x ^ y) == 1) return;
 Basis res = *this;
                                                                add_edge(x ^ 1, y); add_edge(y ^ 1, x);
  for (int i = 0; i < BITS; i++)</pre>
   if (res.b[i].second < l) res.b[i] = {0, MIN};</pre>
                                                               vector<int> solve2sat() {
  return res;
                                                                solve(); vector<int> res(n);
                                                                for (int i = 0; i < n; i += 2)
  if (idx[i] == idx[i + 1]) return {};</pre>
 }
};
                                                                for (int i = 0; i < n; i++)</pre>
2.6 Binary Search On Segtree [6c61c0]
                                                                 res[i] = idx[i] < idx[i ^ 1];
// find_first = l -> minimal x s.t. check( [l, x) )
// find_last = r -> maximal x s.t. check( [x, r) )
                                                                return res;
int find_first(int l, auto &&check) {
 if (l >= n) return n + 1;
                                                              3.3 BCC [6ac6db]
 l += sz; push(l); Monoid sum; // identity
                                                              class BCC {
  while ((l & 1) == 0) l >>= 1;
                                                               int n, ecnt, bcnt;
  if (auto s = sum + nd[l]; check(s)) {
                                                               vector<vector<pair<int, int>>> g;
   while (l < sz) {</pre>
                                                               vector<int> dfn, low, bcc, stk;
                                                               vector<bool> ap, bridge;
void dfs(int u, int f) {
    prop(l); l = (l << 1);
    if (auto nxt = sum + nd[l]; not check(nxt))
                                                                dfn[u] = low[u] = dfn[f] + 1;
     sum = nxt, l++;
```

```
int ch = 0;
                                                                  } else {
                                                                  for (int &x = nx[u]; x != -1 &&
  for (auto [v, t] : g[u]) if (bcc[t] == -1) {
   bcc[t] = 0; stk.push_back(t);
                                                                     in[x] \le in[v] \& in[v] \le out[x]; x = nx[x])
   if (dfn[v]) {
                                                                   merge(u, x);
    low[u] = min(low[u], dfn[v]);
                                                                  up[u]--;
    continue;
                                                                 }
   ++ch, dfs(v, u);
                                                                out[u] = dfc;
   low[u] = min(low[u], low[v]);
   if (low[v] > dfn[u]) bridge[t] = true;
                                                               for (int i = 0; i < n; i++)</pre>
   if (low[v] < dfn[u]) continue;</pre>
                                                                if (in[i] == -1) dfs(dfs, i, -1);
                                                               for (int i = 0; i < n; i++)
  if (dsu.anc(i) == i) id[i] = cnt++;</pre>
   ap[u] = true;
   while (not stk.empty()) {
    int o = stk.back(); stk.pop_back();
                                                               vector<vector<int>> comps(cnt);
    bcc[o] = bcnt;
                                                               for (int i = 0; i < n; i++)</pre>
    if (o == t) break;
                                                                comps[id[dsu.anc(i)]].push_back(i);
                                                               return comps;
   bcnt += 1;
                                                              } // test @ yosupo judge
  }
                                                              3.6 DMST [f4317e]
  ap[u] = ap[u] and (ch != 1 or u != f);
                                                              using lld = int64_t;
                                                              struct E { int s, t; lld w; }; // 0-base
public:
                                                              struct PQ {
 BCC(int n_{-}) : n(n_{-}), ecnt(0), bcnt(0), g(n), dfn(n),
                                                               struct P {
    low(n), stk(), ap(n) {}
                                                                lld v; int i;
 void add_edge(int u, int v) {
                                                                bool operator>(const P &b) const { return v > b.v; }
 g[u].emplace_back(v, ecnt);
  g[v].emplace_back(u, ecnt++);
                                                               min_heap<P> pq; lld tag;
                                                               void push(P p) { p.v -= tag; pq.emplace(p); }
 void solve() {
                                                               P top() { P p = pq.top(); p.v += tag; return p; }
  bridge.assign(ecnt, false); bcc.assign(ecnt, -1);
                                                               void join(PQ &b) {
  for (int i = 0; i < n; ++i) if (!dfn[i]) dfs(i, i);</pre>
                                                                if (pq.size() < b.pq.size())</pre>
                                                                 swap(pq, b.pq), swap(tag, b.tag);
 int bcc_id(int x) const { return bcc[x];
                                                                while (!b.pq.empty()) push(b.top()), b.pq.pop();
 bool is_ap(int x) const { return ap[x]; }
 bool is_bridge(int x) const { return bridge[x]; }
                                                              };
};
                                                              vector<int> dmst(const vector<E> &e, int n, int root) {
3.4 Round Square Tree [cf6d74]
                                                               vector<PQ> h(n * 2);
struct RST { // be careful about isolate point
                                                               for (int i = 0; i < int(e.size()); ++i)</pre>
 int n; vector<vector<int>> T;
                                                                h[e[i].t].push({e[i].w, i});
 RST(auto &G) : n(int(G.size())), T(n) {
                                                               vector<int> a(n * 2); iota(all(a), 0);
                                                               vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
auto o = [&](auto Y, int x) -> int {
  vector<int> stk, vis(n), low(n);
auto dfs = [&](auto self, int u, int d) -> void {
                                                                return x==a[x] ? x : a[x] = Y(Y, a[x]); };
   low[u] = vis[u] = d; stk.push_back(u);
   for (int v : G[u]) if (!vis[v]) {
                                                               auto S = [&](int i) { return o(o, e[i].s); };
    self(self, v, d + 1);
if (low[v] == vis[u]) {
                                                                int pc = v[root] = n;
                                                               for (int i = 0; i < n; ++i) if (v[i] == -1)</pre>
     int cnt = int(T.size()); T.emplace_back();
                                                                for (int p = i; v[p]<0 || v[p]==i; p = S(r[p])) {</pre>
     for (int x = -1; x != v; stk.pop_back())
                                                                 if (v[p] == i)
                                                                  for (int q = pc++; p != q; p = S(r[p])) {
      T[cnt].push_back(x = stk.back());
     T[u].push_back(cnt); // T is rooted
                                                                   h[p].tag -= h[p].top().v; h[q].join(h[p]);
    } else low[u] = min(low[u], low[v]);
                                                                    pa[p] = a[p] = q;
   } else low[u] = min(low[u], vis[v]);
                                                                 while (S(h[p].top().i) == p) h[p].pq.pop();
  for (int u = 0; u < n; u++)
                                                                 v[p] = i; r[p] = h[p].top().i;
   if (!vis[u]) dfs(dfs, u, 1);
                                                               vector<int> ans;
 } // T may be forest; after dfs, stk are the roots
}; // test @ 2020 Shanghai K
                                                               for (int i = pc - 1; i >= 0; i--) if (v[i] != n) {
                                                                for (int f = e[r[i]].t; f!=-1 && v[f]!=n; f = pa[f])
3.5 Edge TCC [5a2668]
                                                                 v[f] = n;
vector<vector<int>> ETCC(auto &adj) {
                                                                ans.push_back(r[i]);
 const int n = static_cast<int>(adj.size());
 vector<int> up(n), low(n), in, out, nx, id;
                                                               return ans; // default minimize, returns edgeid array
 in = out = nx = id = vector < int > (n, -1);
 int dfc = 0, cnt = 0; Dsu dsu(n);
                                                              3.7
                                                                     Dominator Tree [ea5b7c]
 auto merge = [&](int u, int v) {
 dsu.join(u, v); up[u] += up[v]; };
                                                              struct Dominator {
                                                               vector<vector<int>> g, r, rdom; int tk;
vector<int> dfn, rev, fa, sdom, dom, val, rp;
 auto dfs = [&](auto self, int u, int p) -> void {
  in[u] = low[u] = dfc++;
  for (int v : adj[u]) if (v != u) {
                                                               Dominator(int n) : g(n), r(n), rdom(n), tk(0) {
   if (v == p) { p = -1; continue; }
                                                                dfn = rev = fa = sdom = dom =
   if (in[v] == -1) {
                                                                 val = rp = vector\langle int \rangle (n, -1); }
    self(self, v, u);
                                                               void add_edge(int x, int y) { g[x].push_back(y); }
    if (nx[v] == -1 && up[v] <= 1) {
                                                               void dfs(int x) {
     up[u] += up[v]; low[u] = min(low[u], low[v]);
                                                                rev[dfn[x] = tk] = x;
                                                                fa[tk] = sdom[tk] = val[tk] = tk; tk++;
     continue;
                                                                for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
    if (up[v] == 0) v = nx[v];
    if (low[u] > low[v])
                                                                 r[dfn[u]].push_back(dfn[x]);
     low[u] = low[v], swap(nx[u], v);
   for (; v != -1; v = nx[v]) merge(u, v);
} else if (in[v] < in[u]) {</pre>
                                                               void merge(int x, int y) { fa[x] = y; }
    low[u] = min(low[u], in[v]); up[u]++;
                                                               int find(int x, int c = 0) {
```

```
if (fa[x] == x) return c ? -1 : x;
                                                             }
  if (int p = find(fa[x], 1); p != -1) {
                                                              3.9 Centroid Decomposition [670cdd]
   if (sdom[val[x]] > sdom[val[fa[x]]])
                                                              class Centroid {
    val[x] = val[fa[x]];
                                                               vector<vector<pair<int, int>>> g; // g[u] = {(v, w)}
   fa[x] = p;
                                                               vector<int> pa, dep, vis, sz, mx;
   return c ? p : val[x];
                                                               vector<vector<int64_t>> Dist;
  } else return c ? fa[x] : val[x];
                                                               vector<int64_t> Sub, Sub2;
                                                               vector<int> Cnt, Cnt2;
 vector<int> build(int s, int n) {
// return the father of each node in dominator tree
                                                               void DfsSz(vector<int> &tmp, int x) {
                                                                vis[x] = true, sz[x] = 1, mx[x] = 0;
  dfs(s); // p[i] = -2 if i is unreachable from s
                                                                for (auto [u, w] : g[x]) if (not vis[u]) {
  for (int i = tk - 1; i >= 0; --i) {
                                                                 DfsSz(tmp, u); sz[x] += sz[u];
   for (int u : r[i])
                                                                 mx[x] = max(mx[x], sz[u]);
    sdom[i] = min(sdom[i], sdom[find(u)]);
   if (i) rdom[sdom[i]].push_back(i);
                                                                tmp.push_back(x);
   for (int u : rdom[i]) {
    int p = find(u);
                                                               void DfsDist(int x, int64_t D = 0) {
    dom[u] = (sdom[p] == i ? i : p);
                                                                Dist[x].push_back(D); vis[x] = true;
                                                                for (auto [u, w] : g[x])
   if (i) merge(i, rp[i]);
                                                                 if (not vis[u]) DfsDist(u, D + w);
  }
 vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];</pre>
                                                               void DfsCen(int x, int D, int p) {
                                                                vector<int> tmp; DfsSz(tmp, x);
                                                                int M = int(tmp.size()), C = -1;
  for (int i = 1; i < tk; ++i)</pre>
                                                                for (int u : tmp)
   p[rev[i]] = rev[dom[i]];
                                                                 if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
  return p;
                                                                for (int u : tmp) vis[u] = false;
 } // test @ yosupo judge
                                                                DfsDist(C);
                                                                for (int u : tmp) vis[u] = false;
                                                                pa[C] = p, vis[C] = true, dep[C] = D;
3.8 Edge Coloring [029763]
                                                                for (auto [u, w] : g[C])
// max(d_u) + 1 edge coloring, time: O(NM)
                                                                 if (not vis[u]) DfsCen(u, D + 1, C);
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
                                                              public:
 for (int i = 0; i <= N; i++)</pre>
                                                               Centroid(int N) : g(N), pa(N), dep(N),
vis(N), sz(N), mx(N), Dist(N),
  for (int j = 0; j <= N; j++)</pre>
    C[i][j] = G[i][j] = 0;
                                                                Sub(N), Sub2(N), Cnt(N), Cnt2(N) {}
                                                               void AddEdge(int u, int v, int w) {
void solve(vector<pair<int, int>> &E, int N) {
                                                                g[u].emplace_back(v, w);
 int X[kN] = {}, a;
                                                                g[v].emplace_back(u, w);
 auto update = [&](int u) {
  for (X[u] = 1; C[u][X[u]]; X[u]++);
                                                               void Build() { DfsCen(0, 0, -1); }
                                                               void Mark(int v) {
 auto color = [&](int u, int v, int c) {
                                                                int x = v, z = -1;
  int p = G[u][v];
                                                                for (int i = dep[v]; i >= 0; --i) {
  G[u][v] = G[v][u] = c;
                                                                 Sub[x] += Dist[v][i], Cnt[x]++;
 C[u][c] = v, C[v][c] = u;
                                                                 if (z != -1)
  C[u][p] = C[v][p] = 0;
                                                                  Sub2[z] += Dist[v][i], Cnt2[z]++;
  if (p) X[u] = X[v] = p;
                                                                 x = pa[z = x];
  else update(u), update(v);
                                                                }
  return p;
 };
                                                               int64_t Query(int v) {
 auto flip = [&](int u, int c1, int c2) {
                                                                int64_t res = 0;
  int p = C[u][c1];
                                                                int x = v, z = -1;
  swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
                                                                for (int i = dep[v]; i >= 0; --i) {
                                                                 res += Sub[x] + 1LL * Cnt[x] * Dist[v][i];
  if (!C[u][c1]) X[u] = c1;
                                                                 if (z != -1)
  if (!C[u][c2]) X[u] = c2;
                                                                  res -= Sub2[z] + 1LL * Cnt2[z] * Dist[v][i];
  return p;
                                                                 x = pa[z = x];
 };
 for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {</pre>
                                                                return res;
                                                               }
  auto [u, v] = E[t];
                                                              }; // pa, dep are centroid tree attributes
  int v0 = v, c = X[u], c0 = c, d;
                                                              3.10
                                                                    Lowbit Decomposition [d1d724]
  vector<pair<int, int>> L; int vst[kN] = {};
  while (!G[u][v0]) {
                                                              class LBD {
   L.emplace_back(v, d = X[v]);
if (!C[v][c]) for (a=L.size()-1;a>=0;a--)
                                                               int timer, chains;
                                                               vector<vector<int>> G;
     c = color(u, L[a].first, c);
                                                               vector<int> tl, tr, chain, head, dep, pa;
   else if (!C[u][d]) for (a=L.size()-1;a>=0;a--)
                                                               // chains : number of chain
     color(u, L[a].first, L[a].second);
                                                               // tl, tr[u] : subtree interval in the seq. of u
   else if (vst[d]) break;
                                                               // head[i] : head of the chain i
   else vst[d] = 1, v = C[u][d];
                                                               // chian[u] : chain id of the chain u is on
                                                               void predfs(int u, int f) {
                                                                dep[u] = dep[pa[u] = f] + 1;
  if (!G[u][v0]) {
   for (; v; v = flip(v, c, d), swap(c, d));
                                                                for (int v : G[u]) if (v != f) {
   if (C[u][c0]) { a = int(L.size()) - 1;
                                                                 predfs(v, u);
                                                                 if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
    while (--a >= 0 && L[a].second != c);
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
                                                                  chain[u] = chain[v];
   } else t--;
                                                                if (chain[u] == 0) chain[u] = ++chains;
```

```
void dfschain(int u, int f) {
                                                              for (int y : D[x]) vis[y] = 0;
  tl[u] = timer++;
  if (head[chain[u]] == -1)
                                                             for (int x : ord) { // c4
   head[chain[u]] = u;
  for (int v : G[u])
   if (v != f and chain[v] == chain[u])
                                                               if (rk[z] > rk[x]) --vis[z];
    dfschain(v, u);
  for (int v : G[u])
   if (v != f and chain[v] != chain[u])
    dfschain(v, u);
  tr[u] = timer;
                                                             private:
public:
                                                              using bits = bitset<maxn>;
 LBD(auto &&G_) : n((int)size(G_)),
                                                              bits popped, G[maxn], ans;
 timer(0), chains(0), G(G_{-}), tl(n), tr(n), chain(n), head(n + 1, -1), dep(n), pa(n)
                                                              size_t deg[maxn], deo[maxn], n;
                                                              void sort_by_degree() {
  { predfs(0, 0); dfschain(0, 0); }
                                                               popped.reset();
 PII get_subtree(int u) { return {tl[u], tr[u]}; }
                                                               for (size_t i = 0; i < n; ++i)</pre>
 vector<PII> get_path(int u, int v) {
                                                                deg[i] = G[i].count();
  vector<PII> res;
  while (chain[u] != chain[v]) {
                                                                size_t mi = maxn, id = 0;
   if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
   swap(u, v);
int s = head[chain[u]];
                                                                  mi = deg[id = j];
   res.emplace_back(tl[s], tl[u] + 1);
                                                                popped[deo[i] = id] = 1;
   u = pa[s]:
                                                                  u = G[i]._Find_next(u))
  if (dep[u] < dep[v]) swap(u, v);</pre>
                                                                  --deg[u];
  res.emplace_back(tl[v], tl[u] + 1);
  return res;
}; // 記得在資結上對點的修改要改成對其 dfs 序的修改
3.11 Virtual Tree [ad5cf5]
vector<pair<int, int>> build(vector<int> vs, int r) {
 vector<pair<int, int>> res;
                                                                return:
 sort(vs.begin(), vs.end(), [](int i, int j) {
 return dfn[i] < dfn[j]; });</pre>
 vector<int> s = {r};
 for (int v : vs) if (v != r) {
  if (int o = lca(v, s.back()); o != s.back()) {
   while (s.size() >= 2) {
    if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
                                                               cur = P & ( ~G[ pivot ] );
    res.emplace_back(s[s.size() - 2], s.back());
                                                               */ // or simply choose first
    s.pop_back();
   if (s.back() != o) {
                                                                 u = cur._Find_next(u)) {
    res.emplace_back(o, s.back());
                                                                if (R[u]) continue;
    s.back() = o;
                                                                R[u] = 1;
   }
                                                                BK(R, P \& G[u], X \& G[u]);
                                                                R[u] = P[u] = 0, X[u] = 1;
  s.push_back(v);
 for (size_t i = 1; i < s.size(); ++i)</pre>
                                                             public:
  res.emplace_back(s[i - 1], s[i]);
                                                              void init(size_t n_) {
 return res; // (x, y): x->y
                                                               n = n_{;}
3.12 Tree Hashing [d6a9f9]
                                                               ans.reset();
vector<int> g[maxn]; llu h[maxn];
llu F(llu z) { // xorshift64star from iwiwi
 z \stackrel{\wedge}{=} z >> 12; z \stackrel{\wedge}{=} z << 25; z \stackrel{\wedge}{=} z >> 27;
 return z * 2685821657736338717LL;
                                                              int solve() {
llu hsah(int u, int f) {
                                                               for (size_t i = 0; i < n; ++i)</pre>
                                                                deg[i] = G[i].count();
llu r = 127; // bigger?
 for (int v : g[u]) if (v != f) r += hsah(v, u);
                                                               bits pob, nob = 0; pob.set();
 return h[u] = F(r);
} // test @ UOJ 763 & yosupo library checker
                                                                size_t v = deo[i];
3.13 Mo's Algorithm on Tree
                                                                bits tmp;
dfs u:
                                                                tmp[v] = 1;
push u
 iterate subtree
                                                                pob[v] = 0, nob[v] = 1;
Let P = LCA(u, v) with St(u) \le St(v)
if (P == u) query[St(u), St(v)]
                                                              }
else query[Ed(u), St(v)], query[St(P), St(P)]
                                                             };
3.14
      Count Cycles [c7e8f2]
// ord = sort by deg decreasing, rk[ord[i]] = i
// D[i] = edge point from rk small to rk big
                                                             struct MaxClique {
for (int x : ord) { // c3
 for (int y : D[x]) vis[y] = 1;
                                                              bits G[kN], cs[kN];
 for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
```

```
for (int y : D[x]) for (int z : adj[y])
  if (rk[z] > rk[x]) c4 += vis[z]++;
 for (int y : D[x]) for (int z : adj[y])
} // both are O(M*sqrt(M)), test @ 2022 CCPC guangzhou
      MaximalClique [293730]
// contain a self loop u to u, than u won't in clique
template <size_t maxn> class MaxClique {
  for (size_t i = 0; i < n; ++i) {</pre>
   for (size_t j = 0; j < n; ++j)</pre>
    if (not popped[j] and deg[j] < mi)</pre>
   for (size_t u = G[i]._Find_first(); u < n;</pre>
 void BK(bits R, bits P, bits X) {
  if (R.count() + P.count() <= ans.count()) return;</pre>
  if (not P.count() and not X.count()) {
   if (R.count() > ans.count()) ans = R;
  /* greedily chosse max degree as pivot
bits cur = P | X; size_t pivot = 0, sz = 0;
  for ( size_t u = cur._Find_first() ;
   u < n ; u = cur.\_Find\_next(u))
    if (deg[u] > sz) sz = deg[pivot = u];
  bits cur = P & (~G[(P | X)._Find_first()]);
  for (size_t u = cur._Find_first(); u < n;</pre>
  for (size_t i = 0; i < n; ++i) G[i].reset();</pre>
 void add_edges(int u, bits S) { G[u] = S; }
 void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
  sort_by_degree(); // or simply iota( deo... )
  for (size_t i = n; i < maxn; ++i) pob[i] = 0;</pre>
  for (size_t i = 0; i < n; ++i) {</pre>
   BK(tmp, pob \& G[v], nob \& G[v]);
  return static_cast<int>(ans.count());
3.16 MaximumClique [aee5d8]
constexpr size_t kN = 150; using bits = bitset<kN>;
 int ans, sol[kN], q, cur[kN], d[kN], n;
```

```
void init(int _n) {
 n = _n;
  for (int i = 0; i < n; ++i) G[i].reset();</pre>
void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
 void pre_dfs(vector<int> &v, int i, bits mask) {
  if (i < 4) {
   for (int x : v) d[x] = (int)(G[x] \& mask).count();
  sort(all(v), [&](int x, int y) {
  return d[x] > d[y]; });
 vector<int> c(v.size());
  cs[1].reset(), cs[2].reset();
  int l = max(ans - q + 1, 1), r = 2, tp = 0, k;
  for (int p : v) {
   for (k = 1; (cs[k] & G[p]).any(); ++k);
   if (k >= r) cs[++r].reset();
  cs[k][p] = 1;
  if (k < l) v[tp++] = p;
  for (k = l; k < r; ++k)
   for (auto p = cs[k]._Find_first();
     p < kN; p = cs[k]._Find_next(p))</pre>
    v[tp] = (int)p, c[tp] = k, ++tp;
 dfs(v, c, i + 1, mask);
void dfs(vector<int> &v, vector<int> &c,
   int i, bits mask) {
  while (!v.empty()) {
   int p = v.back(); v.pop_back(); mask[p] = 0;
   if (q + c.back() <= ans) return;</pre>
   cur[q++] = p;
   vector<int> nr;
   for (int x : v) if (G[p][x]) nr.push_back(x);
   if (!nr.empty()) pre_dfs(nr, i, mask & G[p]);
   else if (q > ans) ans = q, copy_n(cur, q, sol);
   c.pop_back(); --q;
 }
int solve() {
 vector<int> v(n); iota(all(v), 0);
  ans = q = 0; pre_dfs(v, 0, bits(string(n, '1')));
  return ans; // sol[0 ~ ans-1]
} cliq; // test @ yosupo judge
```

3.17 Minimum Mean Cycle [e23bc0]

```
// WARNING: TYPE matters
struct Edge { int s, t; llf c; };
llf solve(vector<Edge> &e, int n) {
// O(VE), returns inf if no cycle, mmc otherwise
vector<VI> prv(n + 1, VI(n)), prve = prv;
vector<vector<llf>> d(n + 1, vector<llf>(n, inf));
d[0] = vector<llf>(n, 0);
for (int i = 0; i < n; i++) {
  for (int j = 0; j < (int)e.size(); j++) {</pre>
   auto [s, t, c] = e[j];
   if (d[i][s] < inf && d[i + 1][t] > d[i][s] + c) {
   d[i + 1][t] = d[i][s] + c;
    prv[i + 1][t] = s; prve[i + 1][t] = j;
   }
llf mmc = inf; int st = -1;
 for (int i = 0; i < n; i++) {</pre>
 llf avg = -inf;
  for (int k = 0; k < n; k++) {
   if (d[n][i] < inf - eps)</pre>
   avg = max(avg, (d[n][i] - d[k][i]) / (n - k));
   else avg = inf;
  if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
 if (st == -1) return inf;
vector<int> vst(n), eid, cycle, rho;
 for (int i = n; !vst[st]; st = prv[i--][st]) {
 vst[st]++; eid.emplace_back(prve[i][st]);
  rho.emplace_back(st);
while (vst[st] != 2) {
  int v = rho.back(); rho.pop_back();
 cycle.emplace_back(v); vst[v]++;
```

```
reverse(all(eid)); eid.resize(cycle.size());
  return mmc;
            Flow & Matching
            HopcroftKarp [930040]
struct HK {
   vector<int> l, r, a, p; int ans;
  HK(int n, int m, auto \&g) : l(n,-1), r(m,-1), ans(0) {
    for (bool match = true; match;) {
       match = false; a.assign(n, -1); p.assign(n, -1);
       queue<int> q;
       for (int i = 0; i < n; i++)</pre>
         if (l[i] == -1) q.push(a[i] = p[i] = i);
       // bitset<maxn> nvis, t; nvis.set();
       while (!q.empty()) {
         int z, x = q.front(); q.pop();
         if (l[a[x]] != -1) continue;
         for (int y : g[x]) { // or iterate t = g[x] & nvis
            // nvis.reset(y);
           if (r[y] == -1) {
             for (z = y; z != -1;)
               r[z] = x, swap(l[x], z), x = p[x];
             match = true; ++ans; break;
           } else if (p[r[y]] == -1)
             q.push(z = r[y]), p[z] = x, a[z] = a[x];
      }
    }
  }
             Kuhn Munkres [2c09ed]
struct KM { // maximize, test @ UOJ 80
int n, l, r; lld ans; // fl and fr are the match
  vector<lld> hl, hr; vector<int> fl, fr, pre, q;
  void bfs(const auto &w, int s) {
    vector<int> vl(n), vr(n); vector<lld> slk(n, INF);
    l = r = 0; vr[q[r++] = s] = true;
    const auto check = [\&](int x) \rightarrow bool {
       if (vl[x] || slk[x] > 0) return true;
       vl[x] = true; slk[x] = INF;
       if (fl[x] != -1) return vr[q[r++] = fl[x]] = true;
       while (x != -1) swap(x, fr[fl[x] = pre[x]]);
       return false;
    while (true) {
       while (l < r)
         for (int x = 0, y = q[l++]; x < n; ++x) if (!vl[x])
           if (chmin(slk[x], hl[x] + hr[y] - w[x][y]))
if (pre[x] = y, !check(x)) return;
       lld d = ranges::min(slk);
       for (int x = 0; x < n; ++x)
         vl[x] ? hl[x] += d : slk[x] -= d;
       for (int x = 0; x < n; ++x) if (vr[x]) hr[x] -= d;
       for (int x = 0; x < n; ++x) if (!check(x)) return;
  KM(int n_{, const auto \&w) : n(n_{, ans(0), 
    hl(n), hr(n), fl(n, -1), fr(fl), pre(n), q(n) {
    for (int i = 0; i < n; ++i) hl[i]=ranges::max(w[i]);</pre>
    for (int i = 0; i < n; ++i) bfs(w, i);</pre>
    for (int i = 0; i < n; ++i) ans += w[i][fl[i]];</pre>
```

Flow Models

- · Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}$.

 - For each edge (x, y, l, u), connect x → y with capacity u − l.
 For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \rightarrow v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t\to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f
 eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer. Also, f is a mincost valid flow.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching ${\it M}$ on bipartite graph(X, Y)
 - 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise.

3. $x \in X$ is chosen iff x is unvisited; $y \in Y$ is chosen iff y is visited.

For each edge (x, y, c), connect x → y with (cost, cap) = (c, 1) if c > 0, otherwise connect y → x with (cost, cap) = (-c, 1)
 For each edge with c < 0, sum these cost as K, then increase d(y) by 1,

4. For each vertex v with d(v) > 0, connect $S \rightarrow v$ with (cost, cap) =

5. For each vertex v with d(v) < 0, connect $v \rightarrow T$ with (cost, cap) =

2. DFS from unmatched vertices in X.

1. Consruct super source S and sink T

Minimum cost cyclic flow

decrease d(x) by 1

(0, d(v))

(0, -d(v))

```
6. Flow from S to T, the answer is the cost of the flow C+K
· Maximum density induced subgraph
  1. Binary search on answer, suppose we're checking answer {\cal T}
     Construct a max flow model, let K be the sum of all weights
  3. Connect source s \to v, v \in G with capacity K
  4. For each edge (u,v,w) in G, connect u\to v and v\to u with capacity w 5. For v\in G, connect it with sink v\to t with capacity K+2T-
      \left(\sum_{e \in E(v)} w(e)\right) - 2w(v)
  6. \hat{T} is a valid answer if the maximum flow f < K|V|

    Minimum weight edge cover

  1. For each v \in V create a copy v', and connect u' \to v' with weight
     w(u,v).
  2. Connect v \to v' with weight 2\mu(v), where \mu(v) is the cost of the cheap-
      est edge incident to v.
  3. Find the minimum weight perfect matching on G^\prime
• Project selection cheat sheet: S,T 分別代表 0,1 側,最小化總花費。
    i 為 O 時花費 c
                                      (i, T, c)
    i 為 1 時花費 c
                                      (S, i, c)
    i \in I 有任何一個為 0 時花費 c
                                      (i, w, \infty), (w, T, c)
    i\in I 有任何一個為 1 時花費 c
                                       S, w, c), (w, i, \infty)
    i 為 0 時得到 c
                                      直接得到 c; (S, i, c)
    i 為 1 時得到 c
                                      直接得到 c; (i, T, c)
    i 為 0 , i 為 1 時花費 c
                                      (i, j, c)
    i,j 不同時花費 c
                                      (i, j, c), (j, i, c)
直接得到 c; (S, w, c), (w, i, \infty), (w, j, \infty)
    i,j 同時是 0 時得到 c
    i,j 同時是 1 時得到 c
                                     直接得到 c; (i, w, \infty), (j, w, \infty), (w, T, c)
• Submodular functions minimization 

• For a function f: 2^V \to \mathbb{R}, f is a submodular function iff 

• \forall S, T \subseteq V, f(S) + f(T) \ge f(S \cup T) + f(S \cap T), or 

• \forall X \subseteq Y \subseteq V, x \notin Y, f(X \cup \{x\}) - f(X) \ge f(Y \cup \{x\}) - f(Y). 

• To minimize \sum_{i \in \mathcal{F}} \phi_i(x_i) + \sum_{i < j} \phi_{ij}(x_i, x_j) + \sum_{i < j < k} \psi_{ijk}(x_i, x_j, x_k)
  - If \theta_i(1) \geq \theta_i(0), add edge (S, i, \theta_i(1) - \theta_i(0)) and \theta_i(0) to answer; other-
  wise, (i, T, \theta_i(0) - \theta_i(1)) and \theta_i(1).

Add edges (i, j, \phi_{ij}(0, 1) + \phi_{ij}(1, 0) - \phi_{ij}(0, 0) - \phi_{ij}(1, 1)).

Denote x_{ijk} as helper nodes. Let P = \psi_{ijk}(0, 0, 0) + \psi_{ijk}(0, 1, 1) + \psi_{ijk}(0, 1, 1)
    (j,x_{ijk},P), (k,x_{ijk},P), (x_{ijk},T,P); otherwise (x_{ijk},i,-P), (x_{ijk},j,-P),
     (x_{ijk}, k, -P), (S, x_{ijk}, -P).
    The minimum cut of this graph will be the the minimum value of the
    function above.
4.4
         Dinic [32c53e]
template <typename Cap = int64_t> class Dinic {
 struct E { int to, rev; Cap cap; }; int n, st, ed;
 vector<vector<E>> G; vector<size_t> lv, idx;
 bool BFS(int k) {
  lv.assign(n, 0); idx.assign(n, 0);
   queue<int> bfs; bfs.push(st); lv[st] = 1;
  while (not bfs.empty() and not lv[ed]) {
    int u = bfs.front(); bfs.pop();
    for (auto e: G[u]) if (e.cap >> k and !lv[e.to])
     bfs.push(e.to), lv[e.to] = lv[u] + 1;
   return lv[ed];
 Cap DFS(int u, Cap f = numeric_limits<Cap>::max()) {
  if (u == ed) return f;
  Cap ret = 0;
   for (auto &i = idx[u]; i < G[u].size(); ++i) {</pre>
    auto &[to, rev, cap] = G[u][i];
    if (cap <= 0 or lv[to] != lv[u] + 1) continue;</pre>
    Cap nf = DFS(to, min(f, cap));
ret += nf; cap -= nf; f -= nf;
    G[to][rev].cap += nf;
    if (f == 0) return ret;
   if (ret == 0) lv[u] = 0;
  return ret;
public:
 void init(int n_) { G.assign(n = n_, vector<E>()); }
 void add_edge(int u, int v, Cap c) {
  G[u].push_back({v, int(G[v].size()), c});
  G[v].push_back({u, int(G[u].size())-1, 0});
 Cap max_flow(int st_, int ed_) {
```

```
st = st_, ed = ed_; Cap ret = 0;

for (int i = 63; i >= 0; --i)
   while (BFS(i)) ret += DFS(st);
  return ret;
}; // test @ luogu P3376
4.5 HLPP [198e4e]
template <typename T> struct HLPP {
 struct Edge { int to, rev; T flow, cap; };
 int n, mx; vector<vector<Edge>> adj; vector<T> excess;
 vector<int> d, cnt, active; vector<vector<int>> B;
 void add_edge(int u, int v, int f) {
  Edge a{v, (int)size(adj[v]), 0, f};
Edge b{u, (int)size(adj[u]), 0, 0};
  adj[u].push_back(a), adj[v].push_back(b);
 void enqueue(int v) {
  if (!active[v] && excess[v] > 0 && d[v] < n) {</pre>
   mx = max(mx, d[v]);
   B[d[v]].push_back(v); active[v] = 1;
  }
 void push(int v, Edge &e) {
  T df = min(excess[v], e.cap - e.flow);
  if (df <= 0 || d[v] != d[e.to] + 1) return;</pre>
  e.flow += df, adj[e.to][e.rev].flow -= df;
excess[e.to] += df, excess[v] -= df;
  enqueue(e.to);
 void gap(int k) {
  for (int v = 0; v < n; v++) if (d[v] >= k)
   cnt[d[v]]--, d[v] = n, cnt[d[v]]++;
 void relabel(int v) {
  cnt[d[v]]--; d[v] = n;
  for (auto e : adj[v])
   if (e.cap > e.flow) d[v] = min(d[v], d[e.to] + 1);
  cnt[d[v]]++; enqueue(v);
 void discharge(int v) {
  for (auto &e : adj[v])
   if (excess[v] > 0) push(v, e);
   else break:
  if (excess[v] <= 0) return;</pre>
  if (cnt[d[v]] == 1) gap(d[v]);
  else relabel(v);
 T max_flow(int s, int t) {
  for (auto &e : adj[s]) excess[s] += e.cap;
  cnt[0] = n; enqueue(s); active[t] = 1;
  for (mx = 0; mx >= 0;)
   if (!B[mx].empty()) {
    int v = B[mx].back(); B[mx].pop_back();
    active[v] = 0; discharge(v);
   } else
           --mx;
  return excess[t];
 HLPP(int _n) : n(_n), adj(n), excess(n),
  d(n), cnt(n + 1), active(n), B(n) {}
};
4.6 Global Min-Cut [ae7013]
void add_edge(auto &w, int u, int v, int c) {
 w[u][v] += c; w[v][u] += c; }
auto phase(const auto &w, int n, vector<int> id) {
 vector<lld> g(n); int s = -1, t = -1;
 while (!id.empty()) {
  int c = -1;
  for (int i : id) if (c == -1 || g[i] > g[c]) c = i;
  s = t; t = c;
  id.erase(ranges::find(id, c));
  for (int i : id) g[i] += w[c][i];
 return tuple{s, t, g[t]};
lld mincut(auto w, int n) {
 lld cut = numeric_limits<lld>::max();
 vector<int> id(n); iota(all(id), 0);
 for (int i = 0; i < n - 1; ++i) {
  auto [s, t, gt] = phase(w, n, id);
  id.erase(ranges::find(id, t));
  cut = min(cut, gt);
```

```
for (int j = 0; j < n; ++j)</pre>
                                                                 for (int i = 0; i < int(g[u].size()); ++i) {</pre>
                                                                 auto e = g[u][i]; int v = e.to;
   w[s][j] += w[t][j], w[j][s] += w[j][t];
                                                                  auto nd = d[u] + e.c + h[u] - h[v];
                                                                  if (e.f <= 0 or d[v] <= nd) continue;</pre>
 return cut;
\frac{1}{V} = \frac{1}{V} \left( \frac{V^3}{V^3} \right), can be O(VE + V^2 \log V)?
                                                                 f[v] = {u, i}; up[v] = min(up[u], e.f);
                                                                  q.emplace(-(d[v] = nd), v);
4.7 GomoryHu Tree [5edb29]
vector<tuple<int, int, int>> GomoryHu(int n){
 vector<tuple<int, int, int>> rt;
                                                               if (d[T] == INF_C) return nullopt;
 vector<int> g(n);
                                                               for (size_t i = 0; i < d.size(); ++i) h[i] += d[i];
for (int i = T; i != S; i = f[i].first) {</pre>
 for (int i = 1; i < n; ++i) {</pre>
 int t = g[i];
auto f = flow;
                                                                auto &eg = g[f[i].first][f[i].second];
                                                                eg.f -= up[T]; g[eg.to][eg.r].f += up[T];
  rt.emplace_back(f.max_flow(i, t), i, t);
  f.walk(i); // bfs points that connected to i (use
                                                               return pair{up[T], h[T]};
    edges with .cap > 0)
                                                              7
  for (int j = i + 1; j < n; ++j)</pre>
                                                             public:
   if (g[j]==t&&f.connect(j)) // check if i can reach j
                                                              MCMF(int n) : g(n), f(n), up(n), d(n, INF_C) {}
    g[j] = i;
                                                              void add_edge(int s, int t, F c, C w) {
                                                               g[s].emplace_back(t, int(g[t].size()), c, w);
g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
 return rt;
}
/* for our dinic:
                                                              pair<F, C> solve(int a, int b) {
 * void walk(int) { BFS(0); }
                                                               h.assign(g.size(), 0);
* bool connect(int i) { return lv[i]; } */
                                                               F c = 0; C w = 0;
      Minimum Cost Max Flow [04f9cb]
                                                               while (auto r = step(a, b)) {
template <typename F, typename C> class MCMF {
                                                                c += r->first, w += r->first * r->second;
 static constexpr F INF_F = numeric_limits<F>::max();
                                                                 fill(d.begin(), d.end(), INF_C);
 static constexpr C INF_C = numeric_limits<C>::max();
 struct E { int to, r; F f; C c; };
                                                               return {c, w};
vector<vector<E>>> g; vector<pair<int, int>> f;
vector<bool> inq; vector<F> up; vector<C> d;
                                                             };
 optional<pair<F, C>> step(int S, int T) {
  queue<int> q;
                                                              4.10 Minimum Cost Circulation [3f7d84]
  for (q.push(S), d[S] = 0, up[S] = INF_F;
                                                             template <typename F, typename C>
    not q.empty(); q.pop()) {
                                                             struct MinCostCirculation {
   int u = q.front(); inq[u] = false;
                                                              struct ep { int to; F flow; C cost; };
   if (up[u] == 0) continue;
                                                              int n; vector<int> vis; int visc;
   for (int i = 0; i < int(g[u].size()); ++i) {</pre>
                                                              vector<int> fa, fae; vector<vector<int>> g;
    auto e = g[u][i]; int v = e.to;
                                                              vector<ep> e; vector<C> pi;
    if (e.f <= 0 or d[v] <= d[u] + e.c) continue;</pre>
                                                              MinCostCirculation(int n_) : n(n_), vis(n), visc(0), g
    d[v] = d[u] + e.c; f[v] = \{u, i\};
                                                                  (n), pi(n) {}
    up[v] = min(up[u], e.f);
                                                              void add_edge(int u, int v, F fl, C cs) {
    if (not inq[v]) q.push(v);
                                                               g[u].emplace_back((int)e.size());
    inq[v] = true;
                                                               e.emplace_back(v, fl, cs);
   }
                                                               g[v].emplace_back((int)e.size());
                                                               e.emplace_back(u, 0, -cs);
  if (d[T] == INF_C) return nullopt;
  for (int i = T; i != S; i = f[i].first) {
                                                              C phi(int x) {
   auto &eg = g[f[i].first][f[i].second];
                                                               if (fa[x] == -1) return 0;
   eg.f -= up[T]; g[eg.to][eg.r].f += up[T];
                                                               if (vis[x] == visc) return pi[x];
                                                               vis[x] = visc;
  return pair{up[T], d[T]};
                                                               return pi[x] = phi(fa[x]) - e[fae[x]].cost;
public:
                                                              int lca(int u, int v) {
 MCMF(int n) : g(n),f(n),inq(n),up(n),d(n,INF_C) {}
                                                               for (; u != -1 || v != -1; swap(u, v)) if (u != -1) {
 void add_edge(int s, int t, F c, C w) {
                                                                if (vis[u] == visc) return u;
 g[s].emplace_back(t, int(g[t].size()), c, w);
                                                                vis[u] = visc;
  g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
                                                                u = fa[u];
                                                               }
 pair<F, C> solve(int a, int b) {
                                                               return -1;
 F c = 0; C w = 0;
 while (auto r = step(a, b)) {
                                                              void pushflow(int x, C &cost) {
   c += r->first, w += r->first * r->second;
                                                               int v = e[x ^ 1].to, u = e[x].to;
   ranges::fill(inq, false); ranges::fill(d, INF_C);
                                                                ++visc;
                                                               if (int w = lca(u, v); w == -1) {
  return {c, w};
                                                                while (v != -1)
}
                                                                  swap(x ^= 1, fae[v]), swap(u, fa[v]), swap(u, v);
};
                                                               } else {
4.9 Dijkstra Cost Flow [d0cfd9]
                                                                int z = u, dir = 0; F f = e[x].flow;
                                                                 vector<int> cyc = {x};
template <typename F, typename C> class MCMF {
 static constexpr F INF_F = numeric_limits<F>::max();
                                                                 for (int d : {0, 1})
 static constexpr C INF_C = numeric_limits<C>::max();
                                                                  for (int i = (d ? u : v); i != w; i = fa[i]) {
 struct E { int to, r; F f; C c; };
                                                                  cyc.push_back(fae[i] ^ d);
                                                                   if (chmin(f, e[fae[i] ^ d].flow)) z = i, dir = d;
 vector<vector<E>> g; vector<pair<int, int>> f;
 vector<F> up; vector<C> d, h;
 optional<pair<F, C>> step(int S, int T) {
                                                                for (int i : cyc) {
                                                                 e[i].flow -= f; e[i ^ 1].flow += f;
 priority_queue<pair<C, int>> q;
  q.emplace(d[S] = 0, S), up[S] = INF_F;
                                                                 cost += f * e[i].cost;
  while (not q.empty()) {
   auto [l, u] = q.top(); q.pop();
                                                                 if (dir) x ^= 1, swap(u, v);
   if (up[u] == 0 or l != -d[u]) continue;
                                                                while (u != z)
```

```
swap(x ^= 1, fae[v]), swap(u, fa[v]), swap(u, v);
                                                                  match = st = pa = S = vis = slack;
                                                                 rep(u, 1, n) rep(v, 1, n) g[u][v] = \{u, v, 0\};
  }
 }
 void dfs(int u) {
                                                                int ED(edge e) {
  vis[u] = visc:
                                                                 return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
                                                                void update_slack(int u, int x, int &s) {
  for (int i : g[u])
   if (int v = e[i].to; vis[v] != visc and e[i].flow)
                                                                  if (!s || ED(g[u][x]) < ED(g[s][x])) s = u; }</pre>
    fa[v] = u, fae[v] = i, dfs(v);
                                                                void set_slack(int x) {
                                                                  slack[x] = 0;
 C simplex() {
                                                                  for (int u = 1; u <= n; ++u)</pre>
                                                                   if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
  C cost = 0;
  fa.assign(g.size(), -1); fae.assign(e.size(), -1);
                                                                    update_slack(u, x, slack[x]);
  ++visc; dfs(0);
  for (int fail = 0; fail < ssize(e); )</pre>
                                                                void q_push(int x) {
   for (int i = 0; i < ssize(e); i++)</pre>
                                                                  if (x \le n) q.push(x);
    if (e[i].flow and e[i].cost < phi(e[i ^ 1].to) -</pre>
                                                                  else for (int y : flo[x]) q_push(y);
    phi(e[i].to))
     fail = 0, pushflow(i, cost), ++visc;
                                                                void set_st(int x, int b) {
    else ++fail;
                                                                 st[x] = b;
  return cost;
                                                                 if (x > n) for (int y : flo[x]) set_st(y, b);
};
                                                                vector<int> split_flo(auto &f, int xr) {
                                                                  auto it = find(all(f), xr);
4.11 General Graph Matching [5f2293]
                                                                  if (auto pr = it - f.begin(); pr % 2 == 1)
struct Matching {
                                                                   reverse(1 + all(f)), it = f.end() - pr;
 queue<int> q; int ans, n;
                                                                  auto res = vector(f.begin(), it);
 vector<int> fa, s, v, pre, match;
                                                                  return f.erase(f.begin(), it), res;
 int Find(int u) {
  return u == fa[u] ? u : fa[u] = Find(fa[u]); }
                                                                void set_match(int u, int v) {
 int LCA(int x, int y) {
                                                                  match[u] = g[u][v].v;
  static int tk = 0; tk++; x = Find(x); y = Find(y);
                                                                  if (u <= n) return;</pre>
  for (;; swap(x, y)) if (x != n) {
                                                                  int xr = flo_from[u][g[u][v].u];
   if (v[x] == tk) return x;
                                                                  auto &f = flo[u], z = split_flo(f, xr);
   v[x] = tk;
                                                                  rep(i, 0, int(z.size())-1) set_match(z[i], z[i \land 1]);
   x = Find(pre[match[x]]);
                                                                  set_match(xr, v); f.insert(f.end(), all(z));
  }
                                                                void augment(int u, int v) {
 void Blossom(int x, int y, int l) {
                                                                  for (;;) {
  for (; Find(x) != l; x = pre[y]) {
                                                                   int xnv = st[match[u]]; set_match(u, v);
   pre[x] = y, y = match[x];
                                                                   if (!xnv) return;
   if (s[y] == 1) q.push(y), s[y] = 0;
for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
                                                                   set_match(xnv, st[pa[xnv]]);
                                                                   u = st[pa[xnv]], v = xnv;
  }
 bool Bfs(auto &&g, int r) {
                                                                int lca(int u, int v) {
  static int t = 0; ++t;
  iota(all(fa), 0); ranges::fill(s, -1);
q = queue<int>(); q.push(r); s[r] = 0;
                                                                  for (++t; u || v; swap(u, v)) if (u) {
  for (; !q.empty(); q.pop()) {
                                                                   if (vis[u] == t) return u;
   for (int x = q.front(); int u : g[x])
                                                                   vis[u] = t; u = st[match[u]];
    if (s[u] == -1) {
                                                                   if (u) u = st[pa[u]];
     if (pre[u] = x, s[u] = 1, match[u] == n) {
                                                                 }
      for (int a = u, b = x, last;
  b != n; a = last, b = pre[a])
                                                                  return 0;
       last = match[b], match[b] = a, match[a] = b;
                                                                void add_blossom(int u, int o, int v) {
      return true;
                                                                  int b = int(find(n + 1 + all(st), 0) - begin(st));
                                                                  lab[b] = 0, S[b] = 0; match[b] = match[o];
     q.push(match[u]); s[match[u]] = 0;
                                                                  vector<int> f = {o};
    } else if (!s[u] && Find(u) != Find(x)) {
                                                                  for (int x = u, y; x != o; x = st[pa[y]])
f.pb(x), f.pb(y = st[match[x]]), q_push(y);
     int l = LCA(u, x);
     Blossom(x, u, l); Blossom(u, x, l);
                                                                  reverse(1 + all(f));
    }
                                                                  for (int x = v, y; x != o; x = st[pa[y]])
f.pb(x), f.pb(y = st[match[x]]), q_push(y);
  return false;
                                                                  flo[b] = f; set_st(b, b);
                                                                  for (int x = 1; x <= nx; ++x)
 Matching(auto &&g) : ans(0), n(int(g.size())),
                                                                   g[b][x].w = g[x][b].w = 0;
 fa(n+1), s(n+1), v(n+1), pre(n+1, n), match(n+1, n) {
                                                                  for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;</pre>
  for (int x = 0; x < n; ++x)
                                                                  for (int xs : flo[b]) {
   if (match[x] == n) ans += Bfs(g, x);
                                                                   for (int x = 1; x <= nx; ++x)</pre>
 } // match[x] == n means not matched
                                                                    if (g[b][x].w == 0 || ED(g[xs][x]) < ED(g[b][x]))
}; // test @ yosupo judge
                                                                     g[b][x] = g[xs][x], g[x][b] = g[x][xs];
       Weighted Matching [94ca35]
                                                                   for (int x = 1; x <= n; ++x)
                                                                   if (flo_from[xs][x]) flo_from[b][x] = xs;
#define pb emplace_back
#define rep(i, l, r) for (int i=(l); i<=(r); ++i)
                                                                 }
struct WeightGraph { // 1-based
  static const int inf = INT_MAX;
                                                                 set_slack(b);
                                                                void expand_blossom(int b) {
 struct edge { int u, v, w; }; int n, nx;
 vector<int> lab; vector<vector<edge>> g;
                                                                  for (int x : flo[b]) set_st(x, x);
                                                                  int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
 vector<int> slack, match, st, pa, S, vis;
 vector<vector<int>>> flo, flo_from; queue<int> q;
                                                                  for (int x : split_flo(flo[b], xr)) {
 WeightGraph(int n_) : n(n_-), nx(n * 2), lab(nx + 1), g(nx + 1, vector < edge > (nx + 1), slack(nx + 1),
                                                                   if (xs == -1) { xs = x; continue; }
                                                                   pa[xs] = g[x][xs].u; S[xs] = 1, S[x] = 0;
  flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
                                                                   slack[xs] = 0; set_slack(x); q_push(x); xs = -1;
```

15

 $a \geq c \vee b \geq c$

 $n < 0 \lor a = 0$

 $a \geq c \vee b \geq c$

 $n < 0 \lor a = 0$

otherwise

```
n |1234567
                                                                                                                  9
                                                                                                           8
                                                                                                                        10
                                                                                                                              11 12 13 14
                                                                              \binom{2n}{n} 2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8
  for (int x : flo[b])
                                                                             if (x == xr) S[x] = 1, pa[x] = pa[b];
   else S[x] = -1, set_slack(x);
                                                                             5.2 Equations
  st[b] = 0;
                                                                             Stirling Number of the First Kind
 bool on_found_edge(const edge &e) {
                                                                             S_1(n,k) counts the number of permutations of n elements with k disjoint
  if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
                                                                             cycles.
                                                                             S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)
   int nu = st[match[v]]; pa[v] = e.u; S[v] = 1;
                                                                             • S_1(n,i) = [x^i] \left(\prod_{i=0}^{n-1} (x+i)\right), use D&Q and taylor shift.
   slack[v] = slack[nu] = 0; S[nu] = 0; q_push(nu);
  } else if (S[v] == 0) {
                                                                             • S_1(i,k) = \frac{i!}{k!} \left[ x^i \right] \left( \sum_{j \ge 1} \frac{x^j}{j} \right)^k
    if (int o = lca(u, v)) add_blossom(u, o, v);
                                                                             Stirling Number of the Second Kind
   else return augment(u, v), augment(v, u), true;
                                                                             S_2(n,k) counts the number of ways to partition a set of n elements into k
                                                                             nonempty sets.
  return false;
                                                                             • S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)
                                                                             • S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}
 bool matching() {
                                                                             \cdot \ S_2(i,k) = rac{i!}{k!} [x^i] \, (e^x - 1)^k Derivatives/Integrals
  ranges::fill(S, -1); ranges::fill(slack, 0);
  q = queue<int>();
                                                                             Integration by parts: \int_a^b f(x)g(x)dx = [F(x)g(x)]_a^b - \int_a^b F(x)g'(x)dx \frac{d}{dx}\sin^{-1}x = \frac{1}{\sqrt{1-x^2}}\frac{d}{dx}\cos^{-1}x = -\frac{1}{\sqrt{1-x^2}}\frac{d}{dx}\tan^{-1}x = \frac{1}{1+x^2}
  for (int x = 1; x <= nx; ++x)</pre>
   if (st[x] == x \&\& !match[x])
    pa[x] = 0, S[x] = 0, q_push(x);
                                                                               \frac{d}{dx}\tan x = 1 + \tan^2 x \int \tan ax = -\frac{\ln|\cos ax|}{a}
  if (q.empty()) return false;
  for (;;) {
                                                                                \int e^{-x^2} = \frac{\sqrt{\pi}}{2} \text{erf}(x) \int x e^{ax} dx = \frac{e^{ax}}{a^2} (ax - 1)
   while (q.size()) {
     int u = q.front(); q.pop();
                                                                             Extended Euler
     if (S[st[u]] == 1) continue;
                                                                             a^b \equiv \begin{cases} a^{(b \mod \varphi(m)) + \varphi(m)} & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod{m}
     for (int v = 1; v <= n; ++v)</pre>
      if (g[u][v].w > 0 && st[u] != st[v]) {
                                                                             Pentagonal Number Theorem
        if (ED(g[u][v]) != 0)
                                                                             \prod_{n=1}^{\infty} (1 - x^n) = \sum_{k=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2} = (\sum p(n)x^n)^{-1}
        update_slack(u, st[v], slack[st[v]]);
        else if (on_found_edge(g[u][v])) return true;
                                                                             5.3 Extended FloorSum
                                                                               g(a, b, c, n) = \sum_{i=0}^{n} i \lfloor \frac{ai + b}{c} \rfloor
    int d = inf;
    for (int b = n + 1; b <= nx; ++b)</pre>
                                                                                              \left( \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \right)
     if (st[b] == b && S[b] == 1)
                                                                                               +g(a \bmod c, b \bmod c, c, n),
      d = min(d, lab[b] / 2);
    for (int x = 1; x <= nx; ++x)
     if (int s = slack[x]; st[x] == x && s && S[x] <= 0)</pre>
                                                                                               \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
                                                                                              -h(c, c-b-1, a, m-1)),
      d = min(d, ED(g[s][x]) / (S[x] + 2));
    for (int u = 1; u <= n; ++u)</pre>
                                                                               h(a,b,c,n) = \sum_{i=0}^{n} \left\lfloor \frac{ai+b}{c} \right\rfloor^{2}
     if (S[st[u]] == 1) lab[u] += d;
     else if (S[st[u]] == 0) {
      if (lab[u] <= d) return false;</pre>
                                                                                               \left( \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right)
      lab[u] -= d;
                                                                                              +\lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1)
                                                                                              +h(a \bmod c, b \bmod c, c, n)
    rep(b, n + 1, nx) if (st[b] == b \&\& S[b] >= 0)
                                                                                              +2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)
     lab[b] += d * (2 - 4 * S[b]);
                                                                                              +2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
    for (int x = 1; x \le nx; ++x)
                                                                                              0.
     if (int s = slack[x]; st[x] == x &&
                                                                                               nm(m+1) - 2g(c, c-b-1, a, m-1)
       s \&\& st[s] != x \&\& ED(g[s][x]) == 0)
                                                                                              -2f(c,c-b-1,a,m-1)-f(a,b,c,n), otherwise
      if (on_found_edge(g[s][x])) return true;
    for (int b = n + 1; b <= nx; ++b)
                                                                             5.4 Integer Division [cd017d]
     if (st[b] == b && S[b] == 1 && lab[b] == 0)
                                                                             lld fdiv(lld a, lld b)
{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
      expand blossom(b):
                                                                             lld cdiv(lld a, lld b)
  return false;
                                                                              { return a / b + (a % b && (a < 0) ^ (b > 0)); }
 }
 pair<lld, int> solve() {
                                                                             5.5 FloorSum [fb5917]
  ranges::fill(match, 0);
                                                                             // @param n `n < 2^32`
// @param m `1 <= m < 2^32`
  rep(u, 0, n) st[u] = u, flo[u].clear();
  int w_max = 0;
                                                                             // @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
  rep(u, 1, n) rep(v, 1, n) {
                                                                             llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
   flo_from[u][v] = (u == v ? u : 0);
                                                                               llu ans = 0;
   w_max = max(w_max, g[u][v].w);
                                                                               while (true) {
                                                                                if (a >= m) ans += n*(n-1)/2 * (a/m), a %= m;
                                                                                if (b >= m) ans += n * (b/m), b %= m;
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
                                                                                if (llu y_max = a * n + b; y_max >= m) {
  int n_matches = 0; lld tot_weight = 0;
  while (matching()) ++n_matches;
                                                                                 n = (llu)(y_max / m), b = (llu)(y_max % m);
  rep(u, 1, n) if (match[u] \&\& match[u] < u)
                                                                                 swap(m, a);
   tot_weight += g[u][match[u]].w;
                                                                                } else break;
  return make_pair(tot_weight, n_matches);
                                                                               return ans;
 void set_edge(int u, int v, int w) {
  g[u][v].w = g[v][u].w = w; }
                                                                             lld floor_sum(lld n, lld m, lld a, lld b) {
                                                                               llu ans = 0;
                                                                               if (a < 0) {
5
      Math
                                                                                llu a2 = (a \% m + m), d = (a2 - a) / m;
     Common Bounds
                                                                                ans -= 1ULL * n * (n - 1) / 2 * d; a = a2;
```

if (b < 0) {

 $|2\,3\,4\,5\,6\,7\,8\,9\,20\,50\,100\,$ n $|100\,1e3\,1e6\,1e9\,1e12\,$ $|1e15\,$ $|1e18\,$

 $\overline{p(n)}$ 2 3 5 7 11 15 22 30 627 2e5 2e8 $\overline{d(i)}$ 12 32 240 1344 6720 26880 103680

```
llu b2 = (b \% m + m), d = (b2 - b) / m;
                                                                if (P == 2 or n == 0) return n;
  ans -= 1ULL * n * d; b = b2;
                                                                auto check = [&](lld x) {
                                                                 return modpow(int(x), (P - 1) / 2, P); };
 }
                                                                if (check(n) != 1) return -1;
mt19937 rnd(7122); lld z = 1, w;
 return ans + floor_sum_unsigned(n, m, a, b);
                                                                while (check(w = (z * z - n + P) % P) != P - 1)
 5.6 ModMin [253e4d]
                                                                 z = rnd() \% P;
 // min{k | l <= ((ak) mod m) <= r}
                                                                const auto M = [P, w](auto &u, auto &v) {
 optional<llu> mod_min(u32 a, u32 m, u32 l, u32 r) {
                                                                 auto [a, b] = u; auto [c, d] = v;
return make_pair((a * c + b * d % P * w) % P,
 if (a == 0) return l ? nullopt : 0;
 if (auto k = llu(l + a - 1) / a; k * a <= r)</pre>
                                                                    (a * d + b * c) % P);
  return k:
  auto b = m / a, c = m % a;
                                                                pair<lld, lld> r(1, 0), e(z, 1);
for (int q = (P + 1) / 2; q; q >>= 1, e = M(e, e))
 if (auto y = mod_min(c, a, a - r % a, a - l % a))
  return (l + *y * c + a - 1) / a + *y * b;
                                                                 if (q & 1) r = M(r, e);
  return nullopt;
                                                                return int(r.first); // sqrt(n) mod P where P is prime
 5.7 Floor Monoid Product [416e89]
/* template <typename T>
T brute(llu a, llu b, llu c, llu n, T U, T R) {
                                                               5.12 FWT [f82550]
                                                               /* or convolution:
                                                                * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
 for (llu i = 1, l = 0; i <= n; i++, res = res * R)
                                                                * and convolution:
  for (llu \ r = (a*i+b)/c; \ l < r; ++l) res = res * U;
                                                                  x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
                                                               void fwt(int x[], int N, bool inv = false) {
 return res;
 } */
                                                                for (int d = 1; d < N; d <<= 1)
 template <typename T>
                                                                 for (int s = 0; s < N; s += d * 2)
 T euclid(llu a, llu b, llu c, llu n, T U, T R) {
                                                                  for (int i = s; i < s + d; i++) {</pre>
                                                                   int j = i + d, ta = x[i], tb = x[j];
 if (!n) return T{};
                                                                   x[i] = add(ta, tb);
x[j] = sub(ta, tb);
 if (b >= c)
   return mpow(U, b / c) * euclid(a, b % c, c, n, U, R);
 if (a >= c)
 return euclid(a % c, b, c, n, U, mpow(U, a / c) * R);
llu m = (u128(a) * n + b) / c;
                                                                if (inv) {
                                                                 const int invn = modinv(N);
 if (!m) return mpow(R, n);
                                                                 for (int i = 0; i < N; i++)</pre>
 return mpow(R, (c - b - 1) / a) * U
                                                                  x[i] = mul(x[i], invn);
   * euclid(c, (c - b - 1) % a, a, m - 1, R, U)
   * mpow(R, n - (u128(c) * m - b - 1) / a);
                                                               }
                                                               5.13 Packed FFT [0a6af5]
 // time complexity is O(log max(a, b, c))
                                                               VL convolution(const VI &a, const VI &b) {
 // UUUU R UUUUU R ... UUU R 共 N 個 R, 最後一個必是 R
                                                                if (a.empty() || b.empty()) return {};
   一直到第 k 個 R 前總共有 (ak+b)/c 個 U
                                                                const int sz = bit_ceil(a.size() + b.size() - 1);
 5.8 ax+by=gcd [d0cbdd]
                                                                // Should be able to handle N <= 10^5, C <= 10^4
 // ax+ny = 1, ax+ny == ax == 1 \ (mod \ n)
                                                                vector<P> v(sz);
 void exgcd(lld x, lld y, lld &g, lld &a, lld &b) {
                                                                for (size_t i = 0; i < a.size(); ++i) v[i].RE(a[i]);</pre>
 if (y == 0) g = x, a = 1, b = 0;
                                                                for (size_t i = 0; i < b.size(); ++i) v[i].IM(b[i]);</pre>
 else exgcd(y, x \% y, g, b, a), b = (x / y) * a;
                                                                fft(v.data(), sz, /*inv=*/false);
                                                                auto rev = v; reverse(1 + all(rev));
                                                                for (int i = 0; i < sz; ++i)
      Chinese Remainder [d69e74]
 // please ensure r_i\in[0,m_i)
                                                                 P A = (v[i] + conj(rev[i])) / P(2, 0);
 bool crt(lld &m1, lld &r1, lld m2, lld r2) {
                                                                 P B = (v[i] - conj(rev[i])) / P(0, 2);
                                                                 v[i] = A * B;
  if (m2 > m1) swap(m1, m2), swap(r1, r2);
  lld g, a, b; exgcd(m1, m2, g, a, b);
if ((r2 - r1) % g != 0) return false;
                                                                VL c(sz); fft(v.data(), sz, /*inv=*/true);
  m2 /= g; lld D = (r2 - r1) / g % <math>m2 * a % m2;
                                                                for (int i = 0; i < sz; ++i) c[i] = roundl(RE(v[i]));</pre>
  r1 += (D < 0 ? D + m2 : D) * m1; m1 *= m2;
                                                                return c;
  assert (r1 >= 0 && r1 < m1);
   return true;
                                                               VI convolution_mod(const VI &a, const VI &b) {
                                                                if (a.empty() || b.empty()) return {};
                                                                const int sz = bit_ceil(a.size() + b.size() - 1);
 5.10 DiscreteLog [86e463]
                                                                vector<P> fa(sz), fb(sz);
 template<typename Int>
                                                                for (size_t i = 0; i < a.size(); ++i)</pre>
 Int BSGS(Int x, Int y, Int M) {
                                                                 fa[i] = P(a[i] & ((1 << 15) - 1), a[i] >> 15);
 // x^? \setminus equiv y \pmod{M}
                                                                for (size_t i = 0; i < b.size(); ++i)
fb[i] = P(b[i] & ((1 << 15) - 1), b[i] >> 15);
 Int t = 1, c = 0, g = 1;
  for (Int M_ = M; M_ > 0; M_ >>= 1) g = g * x % M;
 for (g = gcd(g, M); t % g != 0; ++c) {
                                                                fft(fa.data(), sz); fft(fb.data(), sz);
                                                                auto rfa = fa; reverse(1 + all(rfa));
   if (t == y) return c;
                                                                for (int i = 0; i < sz; ++i) fa[i] *= fb[i];</pre>
  t = t * x % M;
                                                                for (int i = 0; i < sz; ++i) fb[i] *= conj(rfa[i]);</pre>
                                                                fft(fa.data(), sz, true); fft(fb.data(), sz, true);
 if (y % g != 0) return -1;
                                                                vector<int> res(sz);
 t /= g, y /= g, M /= g;
                                                                for (int i = 0; i < sz; ++i) {</pre>
 Int h = 0, gs = 1;
                                                                 lld A = (lld)roundl(RE((fa[i] + fb[i]) / P(2, 0)));
lld C = (lld)roundl(IM((fa[i] - fb[i]) / P(0, 2)));
 for (; h * h < M; ++h) gs = gs * x % M;
 unordered_map<Int, Int> bs;
                                                                 lld B = (lld)roundl(IM(fa[i])); B %= p; C %= p;
  for (Int s = 0; s < h; bs[y] = ++s) y = y * x % M;
                                                                 res[i] = (A + (B << 15) + (C << 30)) % p;
 for (Int s = 0; s < M; s += h) {</pre>
  t = t * gs % M;
                                                                return res;
  if (bs.count(t)) return c + s + h - bs[t];
                                                               } // test @ yosupo judge with long double
 return -1;
                                                               5.14 CRT for arbitrary mod [e4dde7]
                                                               const int mod = 1000000007;
 5.11 Quadratic Residue [f0baec]
                                                               const int M1 = 985661441; // G = 3 for M1, M2, M3
int get_root(int n, int P) { // ensure 0 <= n < p</pre>
                                                               const int M2 = 998244353;
```

```
const int M3 = 1004535809;
                                                                A.insert(A.begin(), 0);
int superBigCRT(lld A, lld B, lld C) {
                                                                fi(1, A.size()) A[i] = mul(modinv(int(i)), A[i]);
  static_assert (M1 < M2 && M2 < M3);</pre>
                                                                return A;
  constexpr lld r12 = modpow(M1, M2-2, M2);
  constexpr lld r13 = modpow(M1, M3-2, M3);
                                                              S Ln(const S &A) { // coef[0] == 1; res[0] == 0
  constexpr lld r23 = modpow(M2, M3-2, M3);
                                                                auto B = Sx(Mul(Dx(A), Inv(A), bit_ceil(A.size()*2)));
  constexpr lld M1M2 = 1LL * M1 * M2 % mod;
                                                                return B.resize(A.size()), B;
 B = (B - A + M2) * r12 % M2;
 C = (C - A + M3) * r13 % M3;
                                                               S Exp(const S &v) { // coef[0] == 0; res[0] == 1
                                                               return Newton(v, 1,
  [](S &X, S &A, int sz) {
 C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
                                                                  auto Y = X; Y.resize(sz / 2); Y = Ln(Y);
                                                                  fi(0, Y.size()) Y[i] = sub(A[i], Y[i]);
5.15 NTT / FFT [41c1f2]
                                                                  Y[0] = add(Y[0], 1); X = Mul(X, Y, sz); );
template <int mod, int G, int maxn> struct NTT {
 static_assert (maxn == (maxn & -maxn));
                                                               S Pow(S a, lld M) { // period mod*(mod-1)
 int roots[maxn];
                                                                assert(!a.empty() && a[0] != 0);
 NTT () {
                                                                const auto imul = [&a](int s) {
  int r = modpow(G, (mod - 1) / maxn);
                                                               for (int &x: a) x = mul(x, s); }; int c = a[0];
imul(modinv(c)); a = Ln(a); imul(int(M % mod));
a = Exp(a); imul(modpow(c, int(M % (mod - 1))));
  for (int i = maxn >> 1; i; i >>= 1) {
   roots[i] = 1;
   for (int j = 1; j < i; j++)
                                                                return a; // mod x^N where N=a.size()
   roots[i + j] = mul(roots[i + j - 1], r);
   r = mul(r, r);
                                                               S Sqrt(const S &v) { // need: QuadraticResidue
   // for (int j = 0; j < i; j++) // FFT (tested)
                                                                assert(!v.empty() && v[0] != 0);
   // roots[i+j] = polar<llf>(1, PI * j / i);
                                                                const int r = get_root(v[0]); assert(r != -1);
  }
                                                                return Newton(v, r,
                                                                 [](S &X, S &A, int sz) {
 // n must be 2^k, and 0 <= F[i] < mod
                                                                  auto Y = X; Y.resize(sz / 2);
 template <typename T>
                                                                  auto B = Mul(A, Inv(Y), sz);
 void operator()(int F[], T n, bool inv = false) {
                                                                  for (int i = 0, inv2 = mod / 2 + 1; i < sz; i++)</pre>
  for (T i = 0, j = 0; i < n; i++) {
                                                                   X[i] = mul(inv2, add(X[i], B[i])); });
   if (i < j) swap(F[i], F[j]);</pre>
   for (T k = n)1; (j^k) < k; k>=1);
                                                               S Mul(auto &&a, auto &&b) {
                                                                const auto n = a.size() + b.size() - 1;
  for (T s = 1; s < n; s *= 2) {
                                                                auto R = Mul(a, b, bit_ceil(n));
   for (T i = 0; i < n; i += s * 2) {
                                                                return R.resize(n), R;
    for (T j = 0; j < s; j++) {
     int a = F[i+j], b = mul(F[i+j+s], roots[s+j]);
                                                               S MulT(S a, S b, size_t k) {
     F[i+j] = add(a, b); // a + b
                                                                assert(b.size()); reverse(all(b)); auto R = Mul(a, b);
     F[i+j+s] = sub(a, b); // a - b
                                                                R = vector(R.begin() + b.size() - 1, R.end());
    }
                                                                return R.resize(k), R;
   }
                                                               S Eval(const S &f, const S &x) {
  if (inv) {
                                                                if (f.empty()) return vector(x.size(), 0);
   int iv = modinv(int(n));
                                                                const int n = int(max(x.size(), f.size()));
   for (T i = 0; i < n; i++) F[i] = mul(F[i], iv);</pre>
                                                                auto q = vector(n \star 2, S(2, 1)); S ans(n);
   reverse(F + 1, F + n);
                                                                fi(0, x.size()) q[i + n][1] = sub(0, x[i]);

for (int i = n - 1; i > 0; i--)
  }
                                                                q[i] = Mul(q[i << 1], q[i << 1 | 1]);
};
                                                                q[1] = MulT(f, Inv(q[1]), n);
5.16
     Formal Power Series [c6b99a]
                                                                for (int i = 1; i < n; i++) {</pre>
                                                                 auto L = q[i << 1], R = q[i << 1 | 1];</pre>
#define fi(l, r) for (size_t i = (l); i < (r); i++)
                                                                q[i << 1 | 0] = MulT(q[i], R, L.size());
q[i << 1 | 1] = MulT(q[i], L, R.size());</pre>
using S = vector<int>;
auto Mul(auto a, auto b, size_t sz) {
 a.resize(sz), b.resize(sz);
 ntt(a.data(), sz); ntt(b.data(), sz);
                                                                for (int i = 0; i < n; i++) ans[i] = q[i + n][0];</pre>
 fi(0, sz) a[i] = mul(a[i], b[i]);
                                                                return ans.resize(x.size()), ans;
 return ntt(a.data(), sz, true), a;
                                                               pair<S, S> DivMod(const S &A, const S &B) {
                                                                assert(!B.empty() && B.back() != 0);
S Newton(const S &v, int init, auto &&iter) {
                                                                if (A.size() < B.size()) return {{}}, A};</pre>
 for (int sz = 2; Q.size() < v.size(); sz *= 2) {</pre>
                                                                const auto sz = A.size() - B.size() + 1;
                                                                S X = B; reverse(all(X)); X.resize(sz);
  S A{begin(v), begin(v) + min(sz, int(v.size()))};
  A.resize(sz * 2), Q.resize(sz * 2);
                                                                S Y = A; reverse(all(Y)); Y.resize(sz);
  iter(Q, A, sz * 2); Q.resize(sz);
                                                                S Q = Mul(Inv(X), Y);
                                                                Q.resize(sz); reverse(all(Q)); X = Mul(Q, B); Y = A;
 }
                                                                fi(0, Y.size()) Y[i] = sub(Y[i], X[i]);
return Q.resize(v.size()), Q;
                                                                while (Y.size() && Y.back() == 0) Y.pop_back();
                                                                while (Q.size() && Q.back() == 0) Q.pop_back();
S Inv(const S &v) { // v[0] != 0
return Newton(v, modinv(v[0]),
[](S &X, S &A, int sz) {
                                                                return {Q, Y};
                                                               } // empty means zero polynomial
                                                               int LinearRecursionKth(S a, S c, int64_t k) {
   ntt(X.data(), sz), ntt(A.data(), sz);
   for (int i = 0; i < sz; i++)
                                                                const auto d = a.size(); assert(c.size() == d + 1);
                                                                const auto sz = bit_ceil(2 * d + 1), o = sz / 2;
    X[i] = mul(X[i], sub(2, mul(X[i], A[i])));
   ntt(X.data(), sz, true); });
                                                                S q = c; for (int &x: q) x = sub(0, x); q[0]=1;
                                                                S p = Mul(a, q); p.resize(sz); q.resize(sz);
for (int r; r = (k & 1), k; k >>= 1) {
S Dx(S A) {
 fi(1, A.size()) A[i - 1] = mul(i, A[i]);
                                                                 fill(d + all(p), 0); fill(d + 1 + all(q), 0);
                                                                ntt(p.data(), sz); ntt(q.data(), sz);
for (size_t i = 0; i < sz; i++)</pre>
 return A.empty() ? A : (A.pop_back(), A);
S Sx(S A) {
                                                                  p[i] = mul(p[i], q[(i + o) & (sz - 1)]);
```

```
for (size_t i = 0, j = 0; j < sz; i++, j++)</pre>
                                                                      return madd(mmul(x, x, n), c, n); };
   q[i] = q[j] = mul(q[i], q[j]);
                                                                     do {
  ntt(p.data(), sz, true); ntt(q.data(), sz, true);
for (size_t i = 0; i < d; i++) p[i] = p[i << 1 | r];
for (size_t i = 0; i <= d; i++) q[i] = q[i << 1];</pre>
                                                                      p = mmul(msub(z = f(f(z)), y = f(y), n), p, n);
                                                                      if (++i &= 63) if (i == (i & -i)) t = gcd(p, n);
                                                                     } while (t == 1);
 } // Bostan-Mori
                                                                     return t == n ? pollard_rho(n) : t;
 return mul(p[0], modinv(q[0]));
                                                                    } // test @ yosupo judge, ~270ms for Q=100 \,
                                                                     // if use montgomery, \sim 70 \text{ms} for Q=100
} // a_n = \sum_{j=0}^{n} a_{n-j}, c_0 \text{ is not used}
5.17 Partition Number [9bb845]
                                                                    5.21 Berlekamp Massey [a94d00]
ans[0] = tmp[0] = 1;
for (int i = 1; i * i <= n; i++) {
                                                                    template <typename T>
                                                                    vector<T> BerlekampMassey(const vector<T> &output) {
 for (int rep = 0; rep < 2; rep++)
for (int j = i; j <= n - i * i; j++)</pre>
                                                                     vector<T> d(output.size() + 1), me, he;
                                                                     for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
modadd(tmp[j], tmp[j-i]);
for (int j = i * i; j <= n; j++)
modadd(ans[j], tmp[j - i * i]);</pre>
                                                                      for (size_t j = 0; j < me.size(); ++j)</pre>
                                                                       d[i] += output[i - j - 2] * me[j];
                                                                      if ((d[i] -= output[i - 1]) == 0) continue;
                                                                      if (me.empty()) {
5.18 Pi Count [715863]
                                                                       me.resize(f = i);
struct S { int rough; lld large; int id; }; lld PrimeCount(lld n) { // n \sim 10^{13} \Rightarrow < 1s
                                                                       continue;
                                                                      vector<T> o(i - f - 1);
T k = -d[i] / d[f]; o.push_back(-k);
 if (n <= 1) return 0;
 const int v = static_cast<int>(sqrtl(n)); int pc = 0;
                                                                      for (T x : he) o.push_back(x * k);
 vector<int> smalls(v + 1), skip(v + 1); vector<S> z;
 for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
                                                                      if (o.size() < me.size()) o.resize(me.size());</pre>
                                                                      for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
 for (int i : views::iota(0, (v + 1) / 2))
z.emplace_back(2*i+1, (n / (2*i+1) + 1) / 2, i);
                                                                      if (i-f+he.size() >= me.size()) he = me, f = i;
                                                                      me = o;
 for (int p = 3; p <= v; ++p)</pre>
  if (smalls[p] > smalls[p - 1]) {
  const int q = p * p; ++pc;
                                                                     return me;
  if (1LL * q * q > n) break;
  skip[p] = 1;
                                                                    5.22 Gauss Elimination [1f5f8c]
  for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
                                                                    using VI = vector<int>;
  int ns = 0;
                                                                    using VVI = vector<VI>;
  for (auto e : z) if (!skip[e.rough]) {
                                                                    pair<VI, VVI> gauss(VVI A, VI b) { // solve Ax=b
   lld d = 1LL * e.rough * p;
                                                                     const int N = (int)A.size(), M = (int)A[0].size();
   e.large += pc - (d <= v ? z[smalls[d] - pc].large :</pre>
                                                                     vector<int> depv, free(M, true); int rk = 0;
    smalls[n / d]);
                                                                     for (int i = 0; i < M; ++i) {</pre>
   e.id = ns; z[ns++] = e;
                                                                      int p = -1;
                                                                      for (int j = rk; j < N; ++j)</pre>
  z.resize(ns);
                                                                       if (p == -1 || abs(A[j][i]) > abs(A[p][i]))
  for (int j = v / p; j >= p; --j) {
  int c = smalls[j] - pc, e = min(j * p + p, v + 1);
                                                                        p = j;
                                                                      if (p == -1 || A[p][i] == 0) continue;
   for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
                                                                      swap(A[p], A[rk]); swap(b[p], b[rk]);
const int inv = modinv(A[rk][i]);
                                                                      for (int &x : A[rk]) x = mul(x, inv);
 lld ans = z[0].large; z.erase(z.begin());
                                                                      b[rk] = mul(b[rk], inv);
for (int j = 0; j < N; ++j) if (j != rk) {</pre>
 for (auto &[rough, large, k] : z) {
  const lld m = n / rough; --k;
                                                                       int z = A[j][i];
  ans -= large - (pc + k);
                                                                       for (int k = 0; k < M; ++k)
  for (auto [p, _, l] : z)
                                                                        A[j][k] = sub(A[j][k], mul(z, A[rk][k]));
   if (l >= k || p * p > m) break;
                                                                       b[j] = sub(b[j], mul(z, b[rk]));
   else ans += smalls[m / p] - (pc + l);
                                                                      depv.push_back(i); free[i] = false; ++rk;
return ans;
  // test @ yosupo library checker w/ n=1e11, 68ms
                                                                     for (int i = rk; i < N; i++)</pre>
5.19 Miller Rabin [fbd812]
                                                                      if (b[i] != 0) return {{}}, {{}}}; // not consistent
bool isprime(llu x) {
                                                                     VI x(M); VVI h;
                                                                     for (int i = 0; i < rk; i++) x[depv[i]] = b[i];
for (int i = 0; i < M; i++) if (free[i]) {</pre>
 auto witn = [&](llu a, int t) {
  for (llu a2; t--; a = a2) {
                                                                      h.emplace_back(M); h.back()[i] = 1;
   a2 = mmul(a, a, x);
   if (a2 == 1 && a != 1 && a != x - 1) return true;
                                                                      for (int j = 0; j < rk; j++)</pre>
                                                                       h.back()[depv[j]] = sub(0, A[j][i]);
  }
                                                                     }
  return a != 1;
                                                                     return \{x, h\}; // solution = x + span(h[i])
 if (x <= 2 || ~x & 1) return x == 2;</pre>
 int t = countr_zero(x-1); llu odd = (x-1) >> t;
                                                                    5.23 Charateristic Polynomial [cd559d]
 for (llu m:
                                                                   #define rep(x, y, z) for (int x=y; x<z; x++)
using VI = vector<int>; using VVI = vector<VI>;
  {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
  if (m % x != 0 && witn(mpow(m % x, odd, x), t))
                                                                    void Hessenberg(VVI &H, int N) {
   return false;
                                                                     for (int i = 0; i < N - 2; ++i) {
  for (int j = i + 1; j < N; ++j) if (H[j][i]) {</pre>
 return true;
} // test @ luogu 143 & yosupo judge, ~1700ms for Q=1e5
                                                                       rep(k, i, N) swap(H[i+1][k], H[j][k]);
// if use montgomery, ~250ms for Q=1e5
                                                                       rep(k, 0, N) swap(H[k][i+1], H[k][j]);
5.20 Pollard Rho [57ad88]
                                                                       break;
// does not work when n is prime or n == 1
                                                                      if (!H[i + 1][i]) continue;
// return any non-trivial factor
llu pollard_rho(llu n) {
                                                                      for (int j = i + 2; j < N; ++j) {</pre>
static mt19937_64 rnd(120821011);
                                                                       int co = mul(modinv(H[i + 1][i]), H[j][i]);
                                                                       rep(k, i, N) subeq(H[j][k], mul(H[i+1][k], co));
 if (!(n & 1)) return 2;
 llu y = 2, z = y, c = rnd() % n, <math>p = 1, i = 0, t;
                                                                       rep(k, 0, N) addeq(H[k][i+1], mul(H[k][j], co));
 auto f = [&](llu x) {
```

```
}
                                                                   return x;
                                                                  }} // use double instead of long double if possible
VI CharacteristicPoly(VVI A) {
                                                                   5.25 Simplex Construction
int N = (int)A.size(); Hessenberg(A, N);
VVI P(N + 1, VI(N + 1)); P[0][0] = 1;
                                                                  Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j for all
                                                                  1 \le j \le m and x_i \ge 0 for all 1 \le \overline{i} \le n.
 for (int i = 1; i <= N; ++i) {</pre>
                                                                   1. In case of minimization, let c_i' = -c_i
  rep(j, 0, i+1) P[i][j] = j ? P[i-1][j-1] : 0;

for (int j = i - 1, val = 1; j >= 0; --j) {
                                                                  2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} A_{ji} x_i \le -b_j
                                                                  3. \sum_{1 \leq i \leq n} A_{ji} x_i = b_j \rightarrow \mathsf{add} \subseteq \mathsf{and} \supseteq.
   int co = mul(val, A[j][i - 1]);
                                                                  4. If x_i has no lower bound, replace x_i with x_i - x_i'
   rep(k, 0, j+1) subeq(P[i][k], mul(P[j][k], co));
                                                                  5.26 Adaptive Simpson [09669e]
   if (j) val = mul(val, A[j][j - 1]);
                                                                  llf simp(llf l, llf r) {
 }
                                                                   llf m = (l + r) / 2;
                                                                   return (f(l) + f(r) + 4.0 * f(m)) * (r - l) / 6.0;
 if (N \& 1) for (int \&x: P[N]) x = sub(0, x);
 return P[N]; // test: 2021 PTZ Korea K
                                                                  llf F(llf L, llf R, llf v, llf eps) {
                                                                   llf M = (L + R) / 2, vl = simp(L, M), vr = simp(M, R);
                                                                   if (abs(vl + vr - v) <= 15 * eps)
5.24 Simplex [c9c93b]
                                                                     return vl + vr + (vl + vr - v) / 15.0;
namespace simplex {
                                                                   return F(L, M, vl, eps / 2.0) +
// maximize c^Tx under Ax \le B and x \ge 0
                                                                  F(M, R, vr, eps / 2.0);
} // call F(l, r, simp(l, r), 1e-6)
5.27 Golden Ratio Search [376bcb]
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
using VD = vector<llf>;
                                                                  llf gss(llf a, llf b, auto &&f) {
using VVD = vector<vector<llf>>;
                                                                     llf r = (sqrt(5)-1)/2, eps = 1e-7;
const llf eps = 1e-9, inf = 1e+9;
int n, m; VVD d; vector<int> p, q;
                                                                     llf x1 = b - r*(b-a), x2 = a + r*(b-a);
                                                                     llf f1 = f(x1), f2 = f(x2);
void pivot(int r, int s) {
 llf inv = 1.0 / d[r][s];
                                                                     while (b-a > eps)
                                                                       if (f1 < f2) { //change to > to find maximum
 for (int i = 0; i < m + 2; ++i)</pre>
                                                                         b = x2; x2 = x1; f2 = f1;
  for (int j = 0; j < n + 2; ++j)
                                                                         x1 = b - r*(b-a); f1 = f(x1);
   if (i != r && j != s)
                                                                       } else {
    d[i][j] -= d[r][j] * d[i][s] * inv;
                                                                         a = x1; x1 = x2; f1 = f2;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;</pre>
 for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
                                                                         x2 = a + r*(b-a); f2 = f(x2);
 d[r][s] = inv; swap(p[r], q[s]);
                                                                     return a;
bool phase(int z) {
                                                                        Geometry
 int x = m + z;
                                                                        Basic Geometry [1d2d70]
 while (true) {
                                                                  6.1
  int s = -1;
                                                                  #define IM imag
  for (int i = 0; i <= n; ++i) {</pre>
                                                                  #define RE real
   if (!z && q[i] == -1) continue;
                                                                  using lld = int64_t;
using llf = long double;
   if (s == -1 || d[x][i] < d[x][s]) s = i;
                                                                  using PT = complex<lld>;
  if (s == -1 \mid | d[x][s] > -eps) return true;
                                                                  using PF = complex<llf>;
  int r = -1;
                                                                  using P = PT;
  for (int i = 0; i < m; ++i) {</pre>
                                                                  llf abs(P p) { return sqrtl(norm(p)); }
                                                                  PF toPF(PT p) { return PF{RE(p), IM(p)}; }
int sgn(lld x) { return (x > 0) - (x < 0); }</pre>
   if (d[i][s] < eps) continue;</pre>
   if (r == -1 ||
                                                                  lld dot(P a, P b) { return RE(conj(a) * b); }
    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
                                                                  lld cross(P a, P b) { return IM(conj(a) * b); }
  if (r == -1) return false;
                                                                  int ori(P a, P b, P c) {
  pivot(r, s);
                                                                   return sgn(cross(b - a, c - a));
                                                                  int quad(P p) {
VD solve(const VVD &a, const VD &b, const VD &c) {
                                                                   return (IM(p) == 0) // use sgn for PF
 m = (int)b.size(), n = (int)c.size();
                                                                     ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
 d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)</pre>
                                                                  int argCmp(P a, P b) {
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
                                                                   // returns 0/+-1, starts from theta = -PI
 p.resize(m), q.resize(n + 1);
                                                                   int qa = quad(a), qb = quad(b);
 for (int i = 0; i < m; ++i)
                                                                   if (qa != qb) return sgn(qa - qb);
 p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
                                                                   return sgn(cross(b, a));
 for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];</pre>
 q[n] = -1, d[m + 1][n] = 1;
                                                                  P rot90(P p) { return P{-IM(p), RE(p)}; }
 int r = 0;
                                                                  template <typename V> llf area(const V & pt) {
 for (int i = 1; i < m; ++i)
                                                                   lld ret = 0;
  if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
                                                                   for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
 if (d[r][n + 1] < -eps) {</pre>
                                                                    ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
  pivot(r, n);
                                                                   return ret / 2.0;
  if (!phase(1) || d[m + 1][n + 1] < -eps)</pre>
   return VD(n, -inf);
                                                                  template <typename V> PF center(const V & pt) {
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
                                                                   P ret = 0; lld A = 0;
                                                                   for (int i = 1; i + 1 < (int)pt.size(); i++) {
    lld cur = cross(pt[i] - pt[0], pt[i+1] - pt[0]);</pre>
   int s = min_element(d[i].begin(), d[i].end() - 1)
       - d[i].begin();
                                                                    ret += (pt[i] + pt[i + 1] + pt[0]) * cur; A += cur;
   pivot(i, s);
  }
                                                                   return toPF(ret) / llf(A * 3);
 if (!phase(0)) return VD(n, inf);
 VD x(n);
                                                                  PF project(PF p, PF q) { // p onto q
 for (int i = 0; i < m; ++i)</pre>
                                                                   return dot(p, q) * q / dot(q, q); // dot<llf>
 if (p[i] < n) x[p[i]] = d[i][n + 1];</pre>
```

```
6.2 2D Convex Hull [ecba37]
// from NaCl, counterclockwise, be careful of n<=2
vector<P> convex_hull(vector<P> v) {
 sort(all(v)); // by X then Y
 if (v[0] == v.back()) return {v[0]};
 int t = 0, s = 1; vector<P> h(v.size() + 1);
 for (int _ = 2; _-
                    --; s = t--, reverse(all(v)))
  for (P p : v) {
   while (t>s && ori(p, h[t-1], h[t-2]) >= 0) t--;
   h[t++] = p;
 return h.resize(t), h;
6.3
      2D Farthest Pair [8b5844]
// p is CCW convex hull w/o colinear points
int n = (int)p.size(), pos = 1; lld ans = 0;
for (int i = 0; i < n; i++) {
  P e = p[(i + 1) % n] - p[i];
  while (cross(e, p[(pos + 1) % n] - p[i]) >
     cross(e, p[pos] - p[i]))
 pos = (pos + 1) % n;
for (int j: {i, (i + 1) % n})
 ans = max(ans, norm(p[pos] - p[j]));
} // tested @ AOJ CGL_4_B
6.4 MinMax Enclosing Rect [e4470c]
// from 8BQube, plz ensure p is strict convex hull
const llf INF = 1e18, qi = acos(-1) / 2 * 3;
pair<llf, llf> solve(const vector<P> &p) {
 llf mx = 0, mn = INF; int n = (int)p.size();
 for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {
#define Z(v) (p[(v) % n] - p[i])
  P = Z(i + 1);
  while (cross(e, Z(u + 1)) > cross(e, Z(u))) ++u;
  while (dot(e, Z(r + 1)) > dot(e, Z(r))) ++r;
  if (!i) l = r + 1;
  while (dot(e, Z(l + 1)) < dot(e, Z(l))) ++l;</pre>
  PD = p[r \% n] - p[l \% n];
  llf H = cross(e, Z(u)) / llf(norm(e));
  mn = min(mn, dot(e, D) * H);
  llf B = sqrt(norm(D)) * sqrt(norm(Z(u)));
  llf deg = (qi - acos(dot(D, Z(u)) / B)) / 2;
  mx = max(mx, B * sin(deg) * sin(deg));
 return {mn, mx};
} // test @ UVA 819
      Minkowski Sum [602806]
// A, B are strict convex hull rotate to min by (X, Y)
vector<P> Minkowski(vector<P> A, vector<P> B) {
 const int N = (int)A.size(), M = (int)B.size();
 vector\langle P \rangle sa(N), sb(M), C(N + M + 1);
 for (int i = 0; i < N; i++) sa[i] = A[(i+1)%N]-A[i];
for (int i = 0; i < M; i++) sb[i] = B[(i+1)%M]-B[i];</pre>
 C[\Theta] = A[\Theta] + B[\Theta];
 for (int i = 0, j = 0; i < N || j < M; ) {</pre>
  P = (j>=M \mid | (i<N \&\& cross(sa[i], sb[j])>=0))
   ? sa[i++] : sb[j++];
 C[i + j] = e;
 partial_sum(all(C), C.begin()); C.pop_back();
 return convex_hull(C); // just to remove colinear
6.6 Segment Intersection [60d016]
struct Seg { // closed segment
 P st, dir; // represent st + t*dir for 0<=t<=1
 Seg(P s, P e) : st(s), dir(e - s) {}
 static bool valid(lld p, lld q) {
    is there t s.t. 0 <= t <= 1 && qt == p ?
  if (q < 0) q = -q, p = -p;
  return 0 <= p && p <= q;
 vector<P> ends() const { return { st, st + dir }; }
template <typename T> bool isInter(T A, P p) {
 if (A.dir == P(0)) return p == A.st; // BE CAREFUL
 return cross(p - A.st, A.dir) == 0 &&
  T::valid(dot(p - A.st, A.dir), norm(A.dir));
template <typename U, typename V>
bool isInter(U A, V B) {
   if (cross(A.dir, B.dir) == 0) { // BE CAREFUL
 bool res = false;
```

```
for (P p: A.ends()) res |= isInter(B, p);
  for (P p: B.ends()) res |= isInter(A, p);
  return res;
 P D = B.st - A.st; lld C = cross(A.dir, B.dir);
 return U::valid(cross(D, B.dir), C) &&
  V::valid(cross(D, A.dir), C);
6.7 Half Plane Intersection [31e216]
struct Line {
P st, ed, dir;
 Line (P s, P e) : st(s), ed(e), dir(e - s) {}
}; using LN = const Line &;
PF intersect(LN A, LN B) {
llf t = cross(B.st - A.st, B.dir) /
  llf(cross(A.dir, B.dir));
 return toPF(A.st) + toPF(A.dir) * t; // C^3 / C^2
bool cov(LN l, LN A, LN B) {
 i128 u = cross(B.st-A.st, B.dir);
 i128 v = cross(A.dir, B.dir);
 // ori(l.st, l.ed, A.st + A.dir*(u/v)) <= 0?
 i128 x = RE(A.dir) * u + RE(A.st - l.st) * v;
 i128 y = IM(A.dir) * u + IM(A.st - l.st) * v;
 return sgn(x*IM(l.dir) - y*RE(l.dir)) * sgn(v) >= 0;
} // x, y are C^3, also sgn<i128> is needed
bool operator<(LN a, LN b) {</pre>
 if (int c = argCmp(a.dir, b.dir)) return c == -1;
 return ori(a.st, a.ed, b.st) < 0;</pre>
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
// the half plane is the LHS when going from st to ed
llf HPI(vector<Line> &q) {
 sort(q.begin(), q.end());
 int n = (int)q.size(), l = 0, r = -1;
 for (int i = 0; i < n; i++) {
  if (i && !argCmp(q[i].dir, q[i-1].dir)) continue;
  while (l < r && cov(q[i], q[r-1], q[r])) --r;</pre>
  while (l < r && cov(q[i], q[l], q[l+1])) ++l;</pre>
  q[++r] = q[i];
 while (l < r && cov(q[l], q[r-1], q[r])) --r;</pre>
 while (l < r && cov(q[r], q[l], q[l+1])) ++l;</pre>
 n = r - l + 1; // q[l .. r] are the lines
 if (n <= 1 || !argCmp(q[l].dir, q[r].dir)) return 0;</pre>
 vector<PF> pt(n);
 for (int i = 0; i < n; i++)</pre>
  pt[i] = intersect(q[i+l], q[(i+1)%n+l]);
 return area(pt);
} // test @ 2020 Nordic NCPC : BigBrother
      SegmentDist (Sausage) [9d8603]
// be careful of abs<complex<int>> (replace _abs below)
llf PointSegDist(P A, Seg B) {
 if (B.dir == P(0)) return _abs(A - B.st);
 if (sgn(dot(A - B.st, B.dir)) >
   sgn(dot(A - B.ed, B.dir)) <= 0)</pre>
  return abs(cross(A - B.st, B.dir)) / _abs(B.dir);
 return min(_abs(A - B.st), _abs(A - B.ed));
llf SegSegDist(const Seg &s1, const Seg &s2) {
 if (isInter(s1, s2)) return 0;
 return min({
   PointSegDist(s1.st, s2),
   PointSegDist(s1.ed, s2),
   PointSegDist(s2.st, s1),
PointSegDist(s2.ed, s1) });
} // test @ QOJ2444 / PTZ19 Summer.D3
6.9 Rotating Sweep Line [8aff27]
struct Event {
 Pd; int u, v;
 bool operator<(const Event &b) const {</pre>
  return sgn(cross(d, b.d)) > 0; }
P makePositive(P z) { return cmpxy(z, 0) ? -z : z; }
void rotatingSweepLine(const vector<P> &p) {
 const int n = int(p.size());
 vector<Event> e; e.reserve(n * (n - 1) / 2);
 for (int i = 0; i < n; i++)</pre>
  for (int j = i + 1; j < n; j++)</pre>
   e.emplace_back(makePositive(p[i] - p[j]), i, j);
 sort(all(e));
```

```
vector<int> ord(n), pos(n);
                                                               evt.emplace_back(q[i], i + N);
iota(all(ord), 0);
                                                              sort(all(evt));
sort(all(ord), [&p](int i, int j) {
                                                              auto vtx = p; sort(all(vtx));
return cmpxy(p[i], p[j]); });
for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
                                                              auto eval = [](const Seg &a, lld x) -> llf {
                                                               if (RE(a.dir) == 0) {
const auto makeReverse = [](auto &v) {
                                                                assert (x == RE(a.st));
                                                                return IM(a.st) + llf(IM(a.dir)) / 2;
 sort(all(v)); v.erase(unique(all(v)), v.end());
 vector<pair<int,int>> segs;
 for (size_t i = 0, j = 0; i < v.size(); i = j) {</pre>
                                                               llf t = (x - RE(a.st)) / llf(RE(a.dir));
  for (; j < v.size() && v[j] - v[i] <= j - i; j++);
segs.emplace_back(v[i], v[j - 1] + 1 + 1);</pre>
                                                               return IM(a.st) + IM(a.dir) * t;
                                                              lld cur_x = 0;
                                                              auto cmp = [&](const Seg &a, const Seg &b) -> bool {
 return segs;
                                                               if (int s = sgn(eval(a, cur_x) - eval(b, cur_x)))
for (size_t i = 0, j = 0; i < e.size(); i = j) {</pre>
                                                                return s == -1;
 /* do here */
                                                               int s = sgn(cross(b.dir, a.dir));
  vector<size_t> tmp;
                                                               if (cur_x != RE(a.st) && cur_x != RE(b.st)) s *= -1;
 for (; j < e.size() && !(e[i] < e[j]); j++)</pre>
                                                               return s == -1;
  tmp.push_back(min(pos[e[j].u], pos[e[j].v]));
 for (auto [l, r] : makeReverse(tmp)) {
                                                              namespace pbds = __gnu_pbds;
                                                              using Tree = pbds::tree<Seg, int, decltype(cmp),</pre>
  reverse(ord.begin() + l, ord.begin() + r);
                                                                 pbds::rb_tree_tag,
   for (int t = l; t < r; t++) pos[ord[t]] = t;</pre>
                                                                 pbds::tree_order_statistics_node_update>;
}
                                                              Tree st(cmp);
                                                              auto answer = [&](P ep) {
                                                               if (binary_search(all(vtx), ep))
6.10
      Polygon Cut [fdd064]
                                                                return 1; // on vertex
using P = PF;
                                                               Seg H(ep, ep); // ??
vector<P> cut(const vector<P>& poly, P s, P e) {
                                                               auto it = st.lower_bound(H);
vector<P> res;
                                                               if (it != st.end() && isInter(it->first, ep))
for (size_t i = 0; i < poly.size(); i++) {</pre>
                                                                return 1; // on edge
 P cur = poly[i], prv = i ? poly[i-1] : poly.back();
                                                               if (it != st.begin() && isInter(prev(it)->first, ep))
 bool side = ori(s, e, cur) < 0;</pre>
                                                                return 1; // on edge
 if (side != (ori(s, e, prv) < 0))
                                                               auto rk = st.order_of_key(H);
   res.push_back(intersect({s, e}, {cur, prv}));
                                                               return rk % 2 == 0 ? 0 : 2; // 0: outside, 2: inside
 if (side)
  res.push_back(cur);
                                                              vector<int> ans(Q);
                                                              for (auto [ep, i] : evt) {
return res;
                                                               cur_x = RE(ep);
                                                               if (i < 0) { // remove
      Point In Simple Polygon [037c52]
                                                                st.erase(edge[~i]);
bool PIP(const vector<P> &p, P z, bool strict = true) {
                                                               } else if (i < N) \{ // insert
 int cnt = 0, n = (int)p.size();
                                                                auto [it, succ] = st.insert({edge[i], i});
for (int i = 0; i < n; i++) {</pre>
                                                                assert (succ);
 P A = p[i], B = p[(i + 1) \% n];
                                                               } else
  if (isInter(Seg(A, B), z)) return !strict;
                                                               ans[i - N] = answer(ep);
 auto zy = IM(z), Ay = IM(A), By = IM(B);
 cnt ^= ((zy<Ay) - (zy<By)) * ori(z, A, B) > 0;
                                                              return ans;
                                                             } // test @ AOJ CGL_3_C
return cnt;
                                                             6.14 Cyclic Ternary Search [162adf]
                                                             int cyclic_ternary_search(int N, auto &&lt_) {
6.12 Point In Hull (Fast) [1c74d8]
                                                              auto lt = [&](int x, int y) {
bool check(P a, P b) {
                                                               return lt_(x % N, y % N); };
return cross(a, b) == 0 && dot(a, b) <= 0;
                                                              int l = 0, r = N; bool up = lt(0, 1);
                                                              while (r - l > 1) {
bool PIH(const vector<P> &h, P z, bool strict = true) {
                                                               int m = (l + r) / 2;
int n = (int)h.size(), a = 1, b = n - 1, r = !strict;
                                                               if (lt(m, 0) ? up : !lt(m, m+1)) r = m;
if (n < 3) return r && check(h[0] - z, h[n-1] - z);</pre>
                                                               else l = m;
if (ori(h[0],h[a],h[b]) > 0) swap(a, b);
 if (ori(h[0],h[a],z) >= r || ori(h[0],h[b],z) <= -r)</pre>
                                                              return (lt(l, r) ? r : l) % N;
  return false;
while (abs(a - b) > 1) {
  int c = (a + b) / 2;
                                                             6.15 Tangent of Points To Hull [8e1343]
  (ori(h[0], h[c], z) > 0 ? b : a) = c;
                                                            pair<int, int> get_tangent(const vector<P> &v, P p) {
                                                              auto gao = [&](int s) {
return ori(h[a], h[b], z) < r;</pre>
                                                               return cyclic_ternary_search(v.size(),
                                                                 [&](int x, int y) {
                                                                  return ori(p, v[x], v[y]) == s; });
     Point In Polygon (Fast) [00590a]
vector<int> PIPfast(vector<P> p, vector<P> q) {
  const int N = int(p.size()), Q = int(q.size());
                                                              }; // test @ codeforces.com/gym/101201/problem/E
                                                              return {gao(1), gao(-1)}; // (a,b):ori(p,v[a],v[b])<0
                                                             } // plz ensure that point strictly out of hull
vector<pair<P, int>> evt;
                                                              // if colinear, returns arbitrary point on line
vector<Seg> edge;
                                                             6.16 Circle Class & Intersection [d5df51]
for (int i = 0; i < N; i++) {</pre>
 int a = i, b = (i + 1) % N;
                                                             llf FMOD(llf x) {
                                                              if (x < -PI) x += PI * 2;
 P A = p[a], B = p[b];
 assert (A < B || B < A); // std::operator<
                                                              if (x > PI) x -= PI * 2;
 if (B < A) swap(A, B);
                                                              return x;
 evt.emplace_back(A, i);
 evt.emplace_back(B, ~i);
                                                            struct Cir { PF o; llf r; };
                                                            // be carefule when tangent
 edge.emplace_back(A, B);
                                                            vector<llf> intersectAngle(Cir a, Cir b) {
for (int i = 0; i < Q; i++)</pre>
                                                             PF dir = b.o - a.o; llf d2 = norm(dir);
```

```
if (norm(a.r - b.r) >= d2) { // <math>norm(x) := |x|^2
                                                                 P z1 = a - b, z2 = a - c; llf D = cross(z1, z2) * 2;
  if (a.r < b.r) return {-PI, PI}; // a in b</pre>
                                                                 auto c1 = dot(a + b, z1), c2 = dot(a + c, z2);
                                                                 P o = rot90(c2 * z1 - c1 * z2) / D;
   else return {}; // b in a
  } else if (norm(a.r + b.r) <= d2) return {};</pre>
                                                                 return { o, abs(o - a) };
 llf dis = abs(dir), theta = arg(dir);
 llf phi = acos((a.r * a.r + d2 - b.r * b.r) /
                                                                Cir minCircleCover(vector<P> p) {
    (2 * a.r * dis)); // is acos_safe needed?
                                                                 assert (!p.empty());
 llf L = FMOD(theta - phi), R = FMOD(theta + phi);
                                                                 ranges::shuffle(p, mt19937(114514));
                                                                 Cir c = { 0, 0 };
for (size_t i = 0; i < p.size(); i++) {</pre>
 return { L, R };
 vector<PF> intersectPoint(Cir a, Cir b) {
                                                                  if (abs(p[i] - c.o) <= c.r) continue;</pre>
 llf d = abs(a.o - b.o);
                                                                  c = { p[i], 0 };
                                                                  for (size_t j = 0; j < i; j++) {
  if (abs(p[j] - c.o) <= c.r) continue;</pre>
  if (d > b.r+a.r || d < abs(b.r-a.r)) return {};</pre>
 llf dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
 PF dir = (a.o - b.o) / d;
                                                                   c.o = (p[i] + p[j]) / llf(2);
 PF u = dir * d1 + b.o;
                                                                   c.r = abs(p[i] - c.o);
 PF v = rot90(dir) * sqrt(max(0.0L, b.r*b.r-d1*d1));
                                                                   for (size_t k = 0; k < j; k++) {</pre>
 return {u + v, u - v};
                                                                    if (abs(p[k] - c.o) <= c.r) continue;</pre>
 } // test @ AOJ CGL probs
                                                                    c = getCircum(p[i], p[j], p[k]);
 6.17 Circle Common Tangent [d97f1c]
                                                                  }
 // be careful of tangent / exact same circle
                                                                 }
 // sign1 = 1 for outer tang, -1 for inter tang
                                                                 return c;
 vector<Line> common_tan(const Cir &a, const Cir &b, int
                                                                } // test @ TIOJ 1093 & luogu P1742
      sign1) {
                                                                      Circle Union [073c1c]
 if (norm(a.o - b.o) < eps) return {};</pre>
 llf d = abs(a.o - b.o), c = (a.r - sign1 * b.r) / d;
                                                                #define eb emplace_back
 PF v = (b.o - a.o) / d;
                                                                struct Teve { // test@SPOJ N=1000, 0.3~0.5s
 if (c * c > 1) return {};
                                                                 PF p; llf a; int add; // point, ang, add
 if (abs(c * c - 1) < eps) {
                                                                 Teve(PF x, llf y, int z) : p(x), a(y), add(z) {}
  PF p = a.o + c * v * a.r;
                                                                 bool operator<(Teve &b) const { return a < b.a; }</pre>
   return {Line(p, p + rot90(b.o - a.o))};
                                                                // strict: x = 0, otherwise x = -1
bool disjunct(Cir &a, Cir &b, int x)
 vector<Line> ret; llf h = sqrt(max(0.0L, 1-c*c));
 for (int sign2 : {1, -1}) {
                                                                { return sgn(abs(a.o - b.o) - a.r - b.r) > x; }
  PF n = c * v + sign2 * h * rot90(v);
                                                                bool contain(Cir &a, Cir &b, int x)
  PF p1 = a.o + n * a.r;
                                                                { return sgn(a.r - b.r - abs(a.o - b.o)) > x; }
  PF p2 = b.o + n * (b.r * sign1);
                                                                vector<llf> CircleUnion(vector<Cir> &c) {
  ret.emplace_back(p1, p2);
                                                                 // area[i] : area covered by at least i circles
                                                                 int N = (int)c.size(); vector<llf> area(N + 1);
 return ret;
                                                                 vector<vector<int>> overlap(N, vector<int>(N));
                                                                 auto g = overlap; // use simple 2darray to speedup
for (int i = 0; i < N; ++i)</pre>
 6.18 Line-Circle Intersection [10786a]
                                                                  for (int j = 0; j < N; ++j) {
  /* c[j] is non-strictly in c[i]. */</pre>
 vector<PF> LineCircleInter(PF p1, PF p2, PF o, llf r) {
 PF ft = p1 + project(o-p1, p2-p1), vec = p2-p1;
 llf dis = abs(o - ft);
                                                                   overlap[i][j] = i != j &&
 if (abs(dis - r) < eps) return {ft};</pre>
                                                                    (sgn(c[i].r - c[j].r) > 0 | |
                                                                      (sgn(c[i].r - c[j].r) == 0 \&\& i < j)) \&\&
 if (dis > r) return {};
 vec = vec * sqrt(r * r - dis * dis) / abs(vec);
return {ft + vec, ft - vec}; // sqrt_safe?
                                                                    contain(c[i], c[j], -1);
                                                                 for (int i = 0; i < N; ++i)</pre>
                                                                  for (int j = 0; j < N; ++j)
g[i][j] = i != j && !(overlap[i][j] ||</pre>
 6.19 Poly-Circle Intersection [8e5133]
 // Divides into multiple triangle, and sum up
                                                                     overlap[j][i] || disjunct(c[i], c[j], -1));
   from 8BQube, test by HDU2892 & AOJ CGL_7_H
                                                                 for (int i = 0; i < N; ++i) {</pre>
 llf _area(PF pa, PF pb, llf r) {
                                                                  vector<Teve> eve; int cnt = 1;
 if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
                                                                  for (int j = 0; j < N; ++j) cnt += overlap[j][i];</pre>
  if (abs(pb) < eps) return 0;</pre>
                                                                  // if (cnt > 1) continue; (if only need area[1])
 llf S, h, theta;
                                                                  for (int j = 0; j < N; ++j) if (g[i][j]) {</pre>
 llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
llf cB = dot(pb, pb-pa) / a / c, B = acos_safe(cB);
                                                                   auto IP = intersectPoint(c[i], c[j]);
                                                                   PF aa = IP[1], bb = IP[0];
 llf cC = dot(pa, pb) / a / b, C = acos_safe(cC);
                                                                   llf A = arg(aa - c[i].o), B = arg(bb - c[i].o);
 if (a > r) {
                                                                   eve.eb(bb, B, 1); eve.eb(aa, A, -1);
   S = (C / 2) * r * r; h = a * b * sin(C) / c;
                                                                   if (B > A) ++cnt;
   if (h < r && B < PI / 2)
   S = (acos\_safe(h/r)*r*r - h*sqrt\_safe(r*r-h*h));
                                                                  if (eve.empty()) area[cnt] += PI*c[i].r*c[i].r;
 } else if (b > r) {
  theta = PI - B - asin_safe(sin(B) / r * a);
                                                                   sort(eve.begin(), eve.end());
   S = 0.5 * a*r*sin(theta) + (C-theta)/2 * r * r;
                                                                   eve.eb(eve[0]); eve.back().a += PI * 2;
 } else
                                                                   for (size_t j = 0; j + 1 < eve.size(); j++) {</pre>
  S = 0.5 * sin(C) * a * b;
                                                                    cnt += eve[j].add;
 return S;
                                                                    area[cnt] += cross(eve[j].p, eve[j+1].p) *.5;
                                                                    llf t = eve[j + 1].a - eve[j].a;
 llf area_poly_circle(const vector<PF> &v, PF 0, llf r)
                                                                    area[cnt] += (t-sin(t)) * c[i].r * c[i].r *.5;
 llf S = 0;
                                                                  }
  for (size_t i = 0, N = v.size(); i < N; ++i)</pre>
                                                                 }
  S += _area(v[i] - 0, v[(i + 1) % N] - 0, r) *
                                                                 return area;
      ori(0, v[i], v[(i + 1) % N]);
 return abs(S);
                                                                6.22 Polygon Union [42e75b]
 6.20 Minimum Covering Circle [92bb15]
                                                                llf polyUnion(const vector<vector<P>> &p) {
                                                                vector<tuple<P, P, int>> seg;
| Cir getCircum(P a, P b, P c)\{ / / P = complex < llf > \}
```

z[f.a][f.b] = z[f.b][f.c] = z[f.c][f.a] = sgn(v);

```
for (int i = 0; i < ssize(p); i++)</pre>
  for (int j = 0, m = int(p[i].size()); j < m; j++)</pre>
                                                                   const auto F = [\&](int x, int y) \{
                                                                    if (z[x][y] > 0 && z[y][x] <= 0)
   seg.emplace_back(p[i][j], p[i][(j + 1) % m], i);
 llf ret = 0; // area of p[i] must be non-negative
                                                                     next.emplace_back(x, y, i);
 for (auto [A, B, i] : seg) {
  vector<pair<llf, int>> evt{{0, 0}, {1, 0}};
for (auto [C, D, j] : seg) {
                                                                   for (const auto &f : now)
                                                                    F(f.a, f.b), F(f.b, f.c), F(f.c, f.a);
   int sc = ori(A, B, C), sd = ori(A, B, D);
                                                                   now = next;
   if (sc != sd && i != j && min(sc, sd) < 0) {</pre>
    llf sa = cross(D-C, A-C), sb = cross(D-C, B-C);
                                                                  return now:
    evt.emplace_back(sa / (sa - sb), sgn(sc - sd));
   } else if (!sc && !sd && j < i
                                                                 // n^2 delaunay: facets with negative z normal of
     && sgn(dot(B - A, D - C)) > 0) {
                                                                 // convexhull of (x, y, x^2 + y^2), use a pseudo-point
    evt.emplace_back(real((C - A) / (B - A)), 1);
                                                                 // (0, 0, inf) to avoid degenerate case
    evt.emplace_back(real((D - A) / (B - A)), -1);
                                                                 // test @ SPOJ CH3D
                                                                 // llf area = 0, vol = 0; // surface area / volume
   }
                                                                 // for (auto [a, b, c]: faces)
                                                                     area += abs(ver(p[a], p[b], p[c]))/2.0,
  for (auto &[q, _] : evt) q = clamp<llf>(q, 0, 1);
  sort(evt.begin(), evt.end());
                                                                     vol += volume(P3(0, 0, 0), p[a], p[b], p[c])/6.0;
  llf sum = 0, last = 0; int cnt = 0;
                                                                 6.25 3D Projection [68f350]
  for (auto [q, c] : evt) {
                                                                 using P3F = valarray<llf>;
   if (!cnt) sum += q - last;
                                                                 P3F toP3F(P3 p) { return {p.x, p.y, p.z}; }
   cnt += c; last = q;
                                                                 llf dot(P3F a, P3F b) {
                                                                  return a[0]*b[0]+a[1]*b[1]+a[2]*b[2];
  ret += cross(A, B) * sum;
                                                                 P3F housev(P3 A, P3 B, int s) {
 return ret / 2;
                                                                  const llf a = abs(A), b = abs(B);
return toP3F(A) / a + s * toP3F(B) / b;
6.23 3D Point [46b73b]
struct P3 {
                                                                 P project(P3 p, P3 q) {
                                                                  P3 o(0, 0, 1);
 lld x, y, z;
                                                                  P3F u = housev(q, o, q.z > 0 ? 1 : -1);
 P3 operator^(const P3 &b) const {
                                                                  auto pf = toP3F(p);
  return {y*b.z-b.y*z, z*b.x-b.z*x, x*b.y-b.x*y};
                                                                  auto np = pf - 2 * u * dot(u, pf) / dot(u, u);
                                                                  return P(np[0], np[1]);
 //Azimuthal angle (longitude) to x-axis. \in [-pi, pi]
                                                                 \} // project p onto the plane q^Tx = 0
 llf phi() const { return atan2(y, x); }
                                                                 6.26 3D Skew Line Nearest Point
 //Zenith angle (latitude) to the z-axis. \in [0, pi]
 llf theta() const { return atan2(sqrt(x*x+y*y),z); }
                                                                 • L_1: \mathbf{v}_1 = \mathbf{p}_1 + t_1 \mathbf{d}_1, L_2: \mathbf{v}_2 = \mathbf{p}_2 + t_2 \mathbf{d}_2
                                                                 • \boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2
P3 ver(P3 a, P3 b, P3 c) { return (b - a) ^ (c - a); }
                                                                 \cdot \ \boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}, \boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}
                                                                 m{\cdot} \ \ m{c}_1 = m{p}_1 + rac{(m{p}_2 - m{p}_1) \cdot m{n}_2}{m{d}_1 \cdot m{n}_2} m{d}_1, m{c}_2 = m{p}_2 + rac{(m{p}_1 - m{p}_2) \cdot m{n}_1}{m{d}_2 \cdot m{n}_1} m{d}_2
lld volume(P3 a, P3 b, P3 c, P3 d) {
return dot(ver(a, b, c), d - a);
                                                                 6.27 Delaunay [3a4ff1]
                                                                 /* please ensure input points are unique */
P3 rotate_around(P3 p, llf angle, P3 axis) {
                                                                 /* A triangulation such that no points will strictly
 llf s = sin(angle), c = cos(angle);
                                                                 inside circumcircle of any triangle. C should be big
P3 u = normalize(axis);
                                                                 enough s.t. the initial triangle contains all points */
 return u*dot(u, p)*(1-c) + p * c + cross(u, p)*s;
                                                                 #define L(i) ((i)==0 ? 2 : (i)-1)
                                                                 #define R(i) ((i)==2 ? 0 : (i)+1)
                                                                 #define F3 for (int i = 0; i < 3; i++)
6.24 3D Convex Hull [01652a]
struct Face {
                                                                 bool is_inf(P z) { return RE(z) \leftarrow -C || RE(z) \rightarrow= C; }
 int a, b, c;
                                                                 bool in_cc(const array<P,3> &p, P q) {
 Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
                                                                  i128 inf_det = 0, det = 0, inf_N, N;
                                                                  F3 {
                                                                   if (is_inf(p[i]) && is_inf(q)) continue;
else if (is_inf(p[i])) inf_N = 1, N = -norm(q);
auto preprocess(const vector<P3> &pt) {
 auto G = pt.begin();
 auto a = find_if(all(pt), [&](P3 z) {
                                                                   else if (is_inf(q)) inf_N = -1, N = norm(p[i]);
 return z != *G; }) - G;
                                                                   else inf_N = 0, N = norm(p[i]) - norm(q);
 auto b = find_if(all(pt), [&](P3 z) {
                                                                   lld D = cross(p[R(i)] - q, p[L(i)] - q);
 return ver(*G, pt[a], z) != P3(0, 0, 0); }) - G;
                                                                   inf_det += inf_N * D; det += N * D;
 auto c = find_if(all(pt), [&](P3 z) {
  return volume(*G, pt[a], pt[b], z) != 0; }) - G;
                                                                  return inf_det != 0 ? inf_det > 0 : det > 0;
 vector<size_t> id;
 for (size_t i = 0; i < pt.size(); i++)</pre>
                                                                 P v[maxn];
  if (i != a && i != b && i != c) id.push_back(i);
                                                                 struct Tri;
 return tuple{a, b, c, id};
                                                                 struct E {
                                                                  Tri *t; int side;
// return the faces with pt indexes
                                                                  E(Tri *t_=0, int side_=0) : t(t_), side(side_) {}
// all points coplanar case will WA
                                                                 struct Tri {
vector<Face> convex_hull_3D(const vector<P3> &pt) {
                                                                  array<int,3> p; array<Tri*,3> ch; array<E,3> e;
Tri(int a=0, int b=0, int c=0) : p{a, b, c}, ch{} {}
 const int n = int(pt.size());
 if (n <= 3) return {}; // be careful about edge case</pre>
 vector<Face> now;
                                                                  bool has_chd() const { return ch[0] != nullptr; }
                                                                  bool contains(int q) const {
 vector<vector<int>> z(n, vector<int>(n));
 auto [a, b, c, ord] = preprocess(pt);
                                                                   F3 if (ori(v[p[i]], v[p[R(i)]], v[q]) < 0)
 now.emplace_back(a, b, c); now.emplace_back(c, b, a);
                                                                    return false;
 for (auto i : ord) {
                                                                   return true;
  vector<Face> next;
  for (const auto &f : now) {
                                                                  bool check(int q) const {
   lld v = volume(pt[f.a], pt[f.b], pt[f.c], pt[i]);
                                                                   return in_cc({v[p[0]], v[p[1]], v[p[2]]}, v[q]); }
   if (v <= 0) next.push_back(f);</pre>
                                                                  pool[maxn * 10], *it, *root;
```

void link(const E &a, const E &b) {

```
if (a.t) a.t->e[a.side] = b;
                                                                    o.y1 = min(o.y1, s->y1); o.y2 = max(o.y2, s->y2);
if (b.t) b.t->e[b.side] = a;
                                                                  return tree+M;
void flip(Tri *A, int a) {
auto [B, b] = A->e[a]; /* flip edge between A,B */
                                                                 bool touch(int x, int y, lld d2, Node *r){
if (!B || !A->check(B->p[b])) return;
                                                                  lld d = (lld) sqrt(d2) +1;
Tri *X = new (it++) Tri(A->p[R(a)], B->p[b], A->p[a]);
                                                                  return x >= r->x1 - d && x <= r->x2 + d &&
Tri *Y = new (it++) Tri(B->p[R(b)], A->p[a], B->p[b]);
                                                                          y >= r->y1 - d \&\& y <= r->y2 + d;
link(E(X, 0), E(Y, 0));
link(E(X, 1), A\rightarrow e[L(a)]); link(E(X, 2), B\rightarrow e[R(b)]); link(E(Y, 1), B\rightarrow e[L(b)]); link(E(Y, 2), A\rightarrow e[R(a)]);
                                                                 using P = pair<lld, int>;
                                                                 void dfs(int x, int y, P &mn, Node *r) {
                                                                  if (!r || !touch(x, y, mn.first, r)) return;
mn = min(mn, P(dis2(r->x, r->y, x, y), r->id));
A->ch = B->ch = {X, Y, nullptr};
flip(X, 1); flip(X, 2); flip(Y, 1); flip(Y, 2);
                                                                  if (r->f == 1 ? y < r->y : x < r->x)
                                                                   dfs(x, y, mn, r\rightarrow L), dfs(x, y, mn, r\rightarrow R);
void add_point(int p) {
Tri *r = root;
                                                                  else
while (r->has_chd()) for (Tri *c: r->ch)
                                                                    dfs(x, y, mn, r\rightarrow R), dfs(x, y, mn, r\rightarrow L);
 if (c && c->contains(p)) { r = c; break; }
array<Tri*, 3> t; /* split into 3 triangles */
                                                                 int query(int x, int y) {
 F3 t[i] = new (it++) Tri(r->p[i], r->p[R(i)], p);
                                                                  P mn(INF, -1); dfs(x, y, mn, root);
F3 link(E(t[i], 0), E(t[R(i)], 1));
                                                                  return mn.second;
F3 link(E(t[i], 2), r->e[L(i)]);
                                                                 }
r->ch = t;
                                                                } tree;
F3 flip(t[i], 2);
                                                                6.30 kd Closest Pair (3D ver.) [84d9eb]
                                                                ilf solve(vector<P> v) {
auto build(const vector<P> &p) {
                                                                 shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,</pre>
it = pool; int n = (int)p.size();
vector<int> ord(n); iota(all(ord), 0);
                                                                  unordered_map<lld, int>>> m;
shuffle(all(ord), mt19937(114514));
root = new (it++) Tri(n, n + 1, n + 2);
                                                                 llf d = dis(v[0], v[1]);
                                                                 auto Idx = [\&d] (llf x) \rightarrow lld {
copy_n(p.data(), n, v); v[n++] = P(-C, -C);
                                                                  return round(x * 2 / d) + 0.1; };
v[n++] = P(C * 2, -C); v[n++] = P(-C, C * 2);
for (int i : ord) add_point(i);
                                                                 auto rebuild_m = [&m, &v, &Idx](int k) {
                                                                  m.clear();
vector<array<int, 3>> res;
                                                                  for (int i = 0; i < k; ++i)</pre>
for (Tri *now = pool; now != it; now++)
                                                                   m[Idx(v[i].x)][Idx(v[i].y)]
  if (!now->has_chd()) res.push_back(now->p);
                                                                     [Idx(v[i].z)] = i;
return res;
                                                                 }; rebuild_m(2);
                                                                 for (size_t i = 2; i < v.size(); ++i) {</pre>
6.28 Build Voronoi [94f000]
                                                                  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
void build_voronoi_cells(auto &&p, auto &&res) {
                                                                      kz = Idx(v[i].z); bool found = false;
vector<vector<int>> adj(p.size());
                                                                  for (int dx = -2; dx \le 2; ++dx) {
 for (auto f: res) F3 {
                                                                    const lld nx = dx + kx;
 int a = f[i], b = f[R(i)];
                                                                    if (m.find(nx) == m.end()) continue;
 if (a >= p.size() || b >= p.size()) continue;
                                                                    auto& mm = m[nx];
                                                                    for (int dy = -2; dy <= 2; ++dy) {
 adj[a].emplace_back(b);
                                                                     const lld ny = dy + ky;
// use `adj` and `p` and HPI to build cells
                                                                     if (mm.find(ny) == mm.end()) continue;
for (size_t i = 0; i < p.size(); i++) {
  vector<Line> ls = frame; // the frame
                                                                     auto& mmm = mm[ny];
                                                                     for (int dz = -2; dz <= 2; ++dz) {
 for (int j : adj[i]) {
                                                                      const lld nz = dz + kz;
  P m = p[i] + p[j], d = rot90(p[j] - p[i]);
                                                                      if (mmm.find(nz) == mmm.end()) continue;
   assert (norm(d) != 0);
                                                                      const int p = mmm[nz];
                                                                      if (dis(v[p], v[i]) < d) {</pre>
   ls.emplace_back(m, m + d); // doubled coordinate
                                                                       d = dis(v[p], v[i]);
 } // HPI(ls)
                                                                       found = true;
                                                                      }
                                                                     }
6.29
       kd Tree (Nearest Point) [f733e5]
struct KDTree {
struct Node {
                                                                  if (found) rebuild_m(i + 1);
 int x, y, x1, y1, x2, y2, id, f; Node *L, *R;
} tree[maxn], *root;
                                                                  else m[kx][ky][kz] = i;
lld dis2(int x1, int y1, int x2, int y2) {
                                                                 return d;
 lld dx = x1 - x2, dy = y1 - y2;
  return dx * dx + dy * dy;
                                                                6.31
                                                                       Simulated Annealing [4e0fe5]
static bool cmpx(Node& a, Node& b) { return a.x<b.x; }</pre>
                                                                llf anneal() {
                                                                 mt19937 rnd_engine(seed);
 static bool cmpy(Node& a, Node& b) { return a.y<b.y; }</pre>
void init(vector<pair<int,int>> &ip) {
                                                                 uniform_real_distribution<llf> rnd(0, 1);
 for (int i = 0; i < ssize(ip); i++)</pre>
                                                                 const llf dT = 0.001;
  tie(tree[i].x, tree[i].y) = ip[i], tree[i].id = i;
                                                                  // Argument p
                                                                 llf S_cur = calc(p), S_best = S_cur;
  root = build(0, (int)ip.size()-1, 0);
                                                                 for (llf T = 2000; T > EPS; T -= dT) {
                                                                  // Modify p to p_prime
const llf S_prime = calc(p_prime);
Node* build(int L, int R, int d) {
 if (L>R) return nullptr;
                                                                  const llf delta_c = S_prime - S_cur;
  int M = (L+R)/2;
 nth_element(tree+L,tree+M,tree+R+1,d%2?cmpy:cmpx);
                                                                  llf prob = min((llf)1, exp(-delta_c / T));
 Node &o = tree[M]; o.f = d \% 2;
                                                                  if (rnd(rnd_engine) <= prob)</pre>
 0.x1 = 0.x2 = 0.x; 0.y1 = 0.y2 = 0.y;
                                                                    S_cur = S_prime, p = p_prime;
 o.L = build(L, M-1, d+1); o.R = build(M+1, R, d+1);
for (Node *s: {o.L, o.R}) if (s) {
                                                                  if (S_prime < S_best) // find min</pre>
                                                                    S_best = S_prime, p_best = p_prime;
  o.x1 = min(o.x1, s->x1); o.x2 = max(o.x2, s->x2);
```

```
if (!rev[i]) { h = 0; continue; }
for (int j = sa[rev[i] - 1]; i + h < n && j + h < n</pre>
 return S_best;
                                                                     && s[i + h] == s[j + h];) ++h;
6.32
      Triangle Centers [adb146]
                                                                   hi[rev[i]] = h ? h-- : 0;
0 = ... // see min circle cover
                                                                  }
G = (A + B + C) / 3;
H = \hat{G} * 3 - 0 * 2; // orthogonal center
                                                                 }
                                                               };
llf a = abs(B - C), b = abs(A - C), c = abs(A - B);
I = (a * A + b * B + c * C) / (a + b + c);
                                                                7.3 Suffix Array Tools [47abd7]
// FermatPoint: minimizes sum of distance
                                                                struct OfflineGetRange : Suffix {
// if max. angle >= 120 deg then vertex
                                                                 vector<vector<pair<int,int>>> qs; int qid;
// otherwise, make eq. triangle AB'C, CA'B, BC'A
// line AA', BB', CC' intersects at P
                                                                 OfflineGetRange(const auto &s)
                                                                  : Suffix(s), qs(n), qid(0) {}
     Stringology
                                                                 int offline_get_range(int x, int len) {
                                                                  return qs[len].emplace_back(rev[x], qid), qid++;
    Hash [ce7fad]
template <int P = 127, int Q = 1051762951>
                                                                 vector<pair<int,int>> solve_get_range() {
class Hash {
                                                                  vector<pair<int,int>> ans(qid); Dsu dsu(n);
vector<int> h, p;
                                                                 for (int i = 1; i < n; i++) qs[hi[i]].eb(i, -1);
for (int i = n - 1; i >= 0; i--)
for (auto [pos, id] : qs[i] | views::reverse)
public:
Hash(const auto &s) : h(s.size()+1), p(s.size()+1) {
  for (size_t i = 0; i < s.size(); ++i)</pre>
                                                                    if (id == -1) dsu.join(pos - 1, pos);
   h[i + 1] = add(mul(h[i], P), s[i]);
                                                                    else ans[id] =
  generate(all(p), [x = 1, y = 1, this]() mutable {
                                                                      {dsu.get_min(pos), dsu.get_max(pos) + 1};
   return y = x, x = mul(x, P), y; });
                                                                  return qs.assign(n), qid = 0, ans;
                                                                 }
 int query(int l, int r) const { // 0-base [l, r)
                                                                };
  return sub(h[r], mul(h[l], p[r - l]));
                                                                template <int LG = 20> struct SparseTableSA : Suffix {
                                                                 array<vector<int>, LG> mn;
};
                                                                 SparseTableSA(const auto &s) : Suffix(s), mn{hi} {
7.2
      Suffix Array [da27c7]
                                                                  for (int l = 0; l + 1 < LG; l++)
auto sais(const auto &s) {
                                                                   for (int i = 0, len = 1 << l; i + len < n; i++)</pre>
 const int n = (int)s.size(), z = ranges::max(s) + 1;
                                                                    mn[l + 1][i] = min(mn[l][i], mn[l][i + len]);
 vector<int> c(z); for (int x : s) ++c[x];
 partial_sum(all(c), begin(c));
                                                                 int lcp(int a, int b) {
 vector<int> sa(n); auto I = views::iota(0, n);
                                                                  if (a == b) return n - a;
 if (ranges::max(c) <= 1) {</pre>
                                                                  a = rev[a]; b = rev[b];
 for (int i : I) sa[--c[s[i]]] = i;
                                                                  if (a > b) swap(a, b);
const int lg = __lg(b - a);
  return sa;
                                                                  return min(mn[lg][a], mn[lg][b - (1 << lg)]);</pre>
 vector<bool> t(n); t[n - 1] = true;
 for (int i = n - 2; i >= 0; --i)
                                                                 int get_range(int x, int len) { // WIP
int a = rev[x] + 1, b = rev[x] + 1;
 t[i] = (s[i]==s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 auto is_lms = views::filter([&t](int x) {
                                                                  for (int l = LG - 1; l >= 0; l--) {
 return x && t[x] && !t[x - 1]; });
                                                                   const int s = 1 << l;
 auto induce = [&] {
                                                                   if (a + s < n && mn[l][a + s] >= len) a += s;
  for (auto x = c; int y : sa)
                                                                   if (b - s >= 0 && mn[l][b - s] >= len) b -= s;
   if (y--) if (!t[y]) sa[x[s[y] - 1]++] = y;
  for (auto x = c; int y : sa | views::reverse)
                                                                  return {b - 1, a};
   if (y--) if (t[y]) sa[--x[s[y]]] = y;
                                                                 7
 };
                                                               };
 vector<int> lms, q(n); lms.reserve(n);
for (auto x = c; int i : I | is_lms) {
                                                                7.4 Ex SAM [58374b]
                                                                struct exSAM {
  q[i] = int(lms.size());
  lms.push_back(sa[--x[s[i]]] = i);
                                                                 int len[maxn * 2], link[maxn * 2]; // maxlen, suflink
                                                                 int next[maxn * 2][maxc], tot; // [0, tot), root = 0
int ord[maxn * 2]; // topo. order (sort by length)
int cnt[maxn * 2]; // occurrence
 induce(); vector<int> ns(lms.size());
 for (int j = -1, nz = 0; int i : sa | is_lms) {
  if (j >= 0) {
                                                                 int newnode() {
   int len = min({n - i, n - j, lms[q[i] + 1] - i});
                                                                  fill_n(next[tot], maxc, 0);
   ns[q[i]] = nz += lexicographical_compare(
                                                                  return len[tot] = cnt[tot] = link[tot] = 0, tot++;
     begin(s) + j, begin(s) + j + len,
begin(s) + i, begin(s) + i + len);
                                                                 void init() { tot = 0, newnode(), link[0] = -1; }
                                                                 int insertSAM(int last, int c) {
                                                                  int cur = next[last][c];
  j = i;
                                                                  len[cur] = len[last] + 1;
 ranges::fill(sa, 0); auto nsa = sais(ns);
                                                                  int p = link[last];
 for (auto x = c; int y : nsa | views::reverse)
                                                                  while (p != -1 && !next[p][c])
 y = lms[y], sa[--x[s[y]]] = y;
                                                                   next[p][c] = cur, p = link[p];
                                                                  if (p == -1) return link[cur] = 0, cur;
 return induce(), sa;
                                                                  int q = next[p][c];
// sa[i]: sa[i]-th suffix is the
                                                                  if (len[p] + 1 == len[q]) return link[cur] = q, cur;
// i-th lexicographically smallest suffix.
                                                                  int clone = newnode();
// hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
                                                                  for (int i = 0; i < maxc; ++i)</pre>
struct Suffix {
                                                                   next[clone][i] = len[next[q][i]] ? next[q][i] : 0;
 int n; vector<int> sa, hi, rev;
                                                                  len[clone] = len[p] + 1;
 Suffix(const auto &s) : n(int(s.size())),
                                                                  while (p != -1 && next[p][c] == q)
 hi(n), rev(n) {
                                                                   next[p][c] = clone, p = link[p];
                                                                  link[link[cur] = clone] = link[q];
  vector<int> _s(n + 1); // _s[n] = 0;
  copy(all(s), begin(_s)); // s shouldn't contain 0
                                                                  link[q] = clone;
                                                                  return cur;
  sa = sais(_s); sa.erase(sa.begin());
  for (int i = 0; i < n; ++i) rev[sa[i]] = i;</pre>
  for (int i = 0, h = 0; i < n; ++i) {
                                                                 void insert(const string &s) {
```

```
for (j = i + 1, k = i; j < n && s[k] <= s[j]; j++)
k = (s[k] < s[j] ? i : k + 1);</pre>
  int cur = 0;
  for (char ch : s) {
   int &nxt = next[cur][int(ch - 'a')];
                                                               // if (i < n / 2 && j >= n / 2) {
                                                               // for min cyclic shift, call duval(s + s)
   if (!nxt) nxt = newnode();
                                                               // then here s.substr(i, n / 2) is min cyclic shift
  cnt[cur = nxt] += 1;
 }
                                                               for (; i <= k; i += j - k)</pre>
                                                                report(i, j - k); // s.substr(l, len)
void build() {
 queue<int> q; q.push(0);
 while (!q.empty()) {
                                                             } // tested @ luogu 6114, 1368 & UVA 719
  int cur = q.front(); q.pop();
                                                             7.9 Main Lorentz [615b8f]
  for (int i = 0; i < maxc; ++i)</pre>
                                                             vector<pair<int, int>> rep[kN]; // 0-base [l, r]
    if (next[cur][i]) q.push(insertSAM(cur, i));
                                                             void main_lorentz(const string &s, int sft = 0) {
                                                              const int n = s.size();
 vector<int> lc(tot);
                                                              if (n == 1) return;
 for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
                                                              const int nu = n / 2, nv = n - nu;
 partial_sum(all(lc), lc.begin());
                                                              const string u = s.substr(0, nu), v = s.substr(nu),
 for (int i = 1; i < tot; ++i) ord[--lc[len[i]]] = i;</pre>
                                                                 ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend());
                                                              main_lorentz(u, sft), main_lorentz(v, sft + nu);
void solve() {
                                                              for (int i = tot - 2; i >= 0; --i)
   cnt[link[ord[i]]] += cnt[ord[i]];
                                                              auto get_z = [](const vector<int> &z, int i) {
                                                               return (0 <= i and i < (int)z.size()) ? z[i] : 0; };</pre>
                                                              auto add_rep = [&](bool left, int c, int l, int k1,
7.5
    KMP [281185]
                                                                 int k2) {
vector<int> kmp(const auto &s) {
                                                               const int L = max(1, l - k2), R = min(l - left, k1);
vector<int> f(s.size());
                                                               if (L > R) return;
for (int i = 1, k = 0; i < (int)s.size(); ++i) {</pre>
                                                               if (left) rep[l].emplace_back(sft + c - R, sft + c -
 while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
                                                                 L);
  f[i] = (k += (s[i] == s[k]));
                                                               else rep[l].emplace_back(sft + c - R - l + 1, sft + c
                                                                   -L-l+1);
return f:
                                                              for (int cntr = 0; cntr < n; cntr++) {</pre>
vector<int> search(const auto &s, const auto &t) {
                                                               int l, k1, k2;
// return 0-indexed occurrence of t in s
                                                               if (cntr < nu) {</pre>
vector<int> f = kmp(t), r;
                                                                l = nu - cntr;
for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
                                                                k1 = get_z(z1, nu - cntr);
 while (k > 0 \&\& s[i] != t[k]) k = f[k - 1];
                                                                k2 = get_z(z2, nv + 1 + cntr);
  k += (s[i] == t[k]);
                                                               } else {
 if (k == (int)t.size()) {
                                                                l = cntr - nu + 1;
   r.push_back(i - t.size() + 1);
                                                                k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
   k = f[k - 1];
                                                                k2 = get_z(z4, (cntr - nu) + 1);
 }
                                                               if (k1 + k2 >= 1)
return r;
                                                                add_rep(cntr < nu, cntr, l, k1, k2);</pre>
7.6 Z value [6a7fd0]
                                                             }
                                                             7.10 BWT [5a9b3a]
vector<int> Zalgo(const string &s) {
vector<int> z(s.size(), s.size());
                                                             vector<int> v[SIGMA];
for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
                                                             void BWT(char *ori, char *res) {
 int j = clamp(r - i, 0, z[i - l]);
                                                              // make ori -> ori + ori
 for (; i + j < z[0] and s[i + j] == s[j]; ++j);
                                                              // then build suffix array
 if (i + (z[i] = j) > r) r = i + z[l = i];
}
                                                             void iBWT(char *ori, char *res) {
  for (int i = 0; i < SIGMA; i++) v[i].clear();</pre>
return z;
                                                              const int len = strlen(ori);
      Manacher [c938a9]
                                                              for (int i = 0; i < len; i++)</pre>
                                                               v[ori[i] - 'a'].push_back(i);
vector<int> manacher(const string &S) {
                                                              vector<int> a;
const int n = (int)S.size(), m = n * 2 + 1;
                                                              for (int i = 0, ptr = 0; i < SIGMA; i++)</pre>
vector<int> z(m);
string t = "."; for (char c: S) t += c, t += '.';
                                                               for (int j : v[i]) {
for (int i = 1, l = 0, r = 0; i < m; ++i) {
  z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
                                                                a.push_back(j);
                                                                ori[ptr++] = 'a' + i;
 while (i - z[i] >= 0 && i + z[i] < m) {</pre>
                                                               }
   if (t[i - z[i]] == t[i + z[i]]) ++z[i];
                                                              for (int i = 0, ptr = 0; i < len; i++) {</pre>
                                                               res[i] = ori[a[ptr]];
   else break;
                                                               ptr = a[ptr];
 if (i + z[i] > r) r = i + z[i], l = i;
                                                              res[len] = 0;
return z; // the palindrome lengths are z[i] - 1
                                                             7.11 Palindromic Tree [0673ee]
/* for (int i = 1; i + 1 < m; ++i) {
                                                             struct PalindromicTree {
 int l = (i - z[i] + 2) / 2, r = (i + z[i]) / 2; if (l != r) / / [l, r) is maximal palindrome
                                                              struct node {
                                                               int nxt[26], f, len; // num = depth of fail link
                                                                                // = #pal_suffix of this node
                                                               int cnt, num;
7.8 Lyndon Factorization [d22cc9]
                                                               node(int l = 0) : nxt{}, f(0), len(l), cnt(0), num(0)
// partition s = w[0] + w[1] + ... + w[k-1],
                                                                  {}
// w[0] >= w[1] >= ... >= w[k-1]
// each w[i] strictly smaller than all its suffix
                                                              vector<node> st; vector<char> s; int last, n;
void duval(const auto &s, auto &&report) {
                                                              void init() {
for (int n = (int)s.size(), i = 0, j, k; i < n; ) {</pre>
                                                              st.clear(); s.clear();
```

```
last = 1; n = 0;
 st.push_back(0); st.push_back(-1);
 st[0].f = 1; s.push_back(-1);
int getFail(int x) {
 while (s[n - st[x].len - 1] != s[n]) x = st[x].f;
  return x;
void add(int c) {
 s.push_back(c -= 'a'); ++n;
 int cur = getFail(last);
  if (!st[cur].nxt[c]) {
  int now = st.size();
   st.push_back(st[cur].len + 2);
  st[now].f = st[getFail(st[cur].f)].nxt[c];
  st[cur].nxt[c] = now;
  st[now].num = st[st[now].f].num + 1;
 last = st[cur].nxt[c]; ++st[last].cnt;
}
void dpcnt() { // cnt = #occurence in whole str
  for (int i = st.size() - 1; i >= 0; i--)
  st[st[i].f].cnt += st[i].cnt;
int size() { return st.size() - 2; }
} pt;
/* usage
string s; cin >> s; pt.init();
for (int i = 0; i < SZ(s); i++) {
int prvsz = pt.size(); pt.add(s[i]);
if (prvsz != pt.size()) {
 int r = i, l = r - pt.st[pt.last].len + 1;
 // pal @ [l,r]: s.substr(l, r-l+1)
} */
```

8 Misc 8.1 **Theorems Spherical Coordinate**





$$r = \sqrt{x^2 + y^2 + z^2}$$

$$\theta = \text{acos}(z/\sqrt{x^2 + y^2 + z^2})$$

$$\phi = \text{atan2}(y, x)$$

Spherical Cap

- · A portion of a sphere cut off by a plane.
- r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ : $\arcsin(a/r)$.
- $\begin{aligned} & \text{Holime} &= \pi h^2 (3r-h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos\theta) (1 \cos\theta)^2/3. \\ & \cdot \text{Area} &= 2\pi r h = \pi (a^2 + h^2) = 2\pi r^2 (1 \cos\theta). \end{aligned}$

Sherman-Morrison formula

$$\left(A + uv^{\mathsf{T}} \right)^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1 + v^{\mathsf{T}}A^{-1}u}$$

Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} =$ d(i), $L_{ij} = -c$ where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is $\det(\tilde{L}_{11})$.
- The number of directed spanning tree rooted at r in G is $\det(\tilde{L}_{rr})$.

BEST Theorem

#{Eulerian circuits} = #{arborescences rooted at 1} $\cdot \prod_{v \in V} (\deg(v) - 1)!$

Random Walk on Graph

Let P be the transition matrix of a strongly connected directed graph, $\sum_{i} P_{i,j} = 1$. Let $F_{i,j}$ be the expected time to reach j from i. Let g_i be the expected time from i to i, G = diag(g) and J be a matrix all of 1, i.e. $J_{i,j} = 1$. Then, F = J - G + PF

First solve G: let $\pi P = \pi$ be a stationary distribution. Then $\pi_i g_i = 1$. The rank of I-P is n-1, so we first solve a special solution X such that (I-P)X: J-G and adjust X to F by $F_{i,j}=X_{i,j}-X_{j,j}$.

Tutte Matrix

For i < j, $d_{ij} = x_{ij}$ (in practice, a random number) if $(i,j) \in E$, otherwise $d_{ij}=0.$ For $i\geq j, d_{ij}=-d_{ji}.$ $rac{{\sf rank}(D)}{2}$ is the maximum matching.

Cayley's Formula

- Given a degree sequence d_1, d_2, \dots, d_n for each labeled vertices, there're $rac{(n-2)!}{-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=$

Erdős-Gallai theorem

A sequence of non-negative integers $d_1 > d_2 > \ldots > d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and $\sum_{i=1}^k d_i \le k(k-1)+\sum_{i=k+1}^n \min(d_i,k)$ holds for all $1 \leq k \leq n$.

Havel-Hakimi algorithm

Find the vertex who has greatest degree unused, connect it with other greatest vertex.

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and b_1, \ldots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{k=1}^k a_i \leq \sum_{i=1}^n \min(b_i, k)$ holds for every $1 \le k \le n$.

Fulkerson-Chen-Anstee theorem

A sequence $(a_1,b_1),\ldots,(a_n,b_n)$ of nonnegative integer pairs with $a_1\geq$ $\cdots \geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq a_i$ $\sum_{i=1}^k \min(b_i, k-1) + \sum_{i=k+1}^n \min(b_i, k)$ holds for every $1 \le k \le n$. **Euler's planar graph formula** V - E + F = C + 1. $E \le 3V - 6$ (when $V \ge 3$)

Pick's theorem

For simple polygon, when points are all integer, we have \boldsymbol{A} #{lattice points in the interior} $+\frac{1}{2}$ #{lattice points on the boundary} -1

Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2)$, find maximum $S\in I_1\cap I_2$. For each iteration, build the directed graph and find a shortest path from s to t.

```
• s \to x : S \sqcup \{x\} \in I_1
x \to t : S \sqcup \{x\} \in I_2
• y \to x : S \setminus \{y\} \sqcup \{x\} \in I_1 (y is in the unique circuit of S \sqcup \{x\})
• x \to y : S \setminus \{y\} \sqcup \{x\} \in I_2 (y is in the unique circuit of S \sqcup \{x\})
```

Alternate the path, and |S| will increase by 1. In each iteration, |E| = O(RN), where $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|.$ For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Find the shortest path by Bellman-Ford. The maximum iteration of Bellman-Ford is 2R + 1.

Dual of LP

Primal	Dual
Maximize $c^{T}x$ s.t. $Ax \leq b$, $x \geq 0$	Minimize $b^{T}y$ s.t. $A^{T}y \geq c$, $y \geq 0$
Maximize $c^{T}x$ s.t. $Ax \leq b$	Minimize $b^{T}y$ s.t. $A^{T}y = c, y \geq 0$
Maximize $c^{T}x$ s.t. $Ax = b, x \geq 0$	Minimize $b^{T}y$ s.t. $A^{T}y \geq c$

Minimax Theorem

Let $f: X \times Y \to \mathbb{R}$ be continuous where $X \subseteq \mathbb{R}^n, Y \subseteq \mathbb{R}^m$ are compact and convex. If $f(\cdot,y):X\to\mathbb{R}$ is concave for fixed y, and $f(x,\cdot):Y\to\mathbb{R}$ is convex for fixed x, then $\max_{x \in X} \min_{y \in Y} f(x, y) = \min_{y \in Y} \max_{x \in X} f(x, y)$, e.g. $f(x, y) = x^{\mathsf{T}} A y$ for

zero-sum matrix game. Parallel Axis Theorem

The second moment of area is $I_z=\iint x^2+y^2\mathrm{d}A.\ I_{z'}=I_z+Ad^2$ where d is the distance between two parallel axis z, z'

8.2 Weight Matroid Intersection [d00ee8]

```
struct Matroid {
Matroid(bitset<N>); // init from an independent set
bool can_add(int); // check if break independence
 Matroid remove(int); // removing from the set
auto matroid_intersection(const vector<int> &w) {
 const int n = (int)w.size(); bitset<N> S;
 for (int sz = 1; sz <= n; sz++) {</pre>
  Matroid M1(S), M2(S); vector<vector<pii>>> e(n + 2);
  for (int j = 0; j < n; j++) if (!S[j]) {
  if (M1.can_add(j)) e[n].eb(j, -w[j]);</pre>
   if (M2.can_add(j)) e[j].eb(n + 1, 0);
  for (int i = 0; i < n; i++) if (S[i]) {</pre>
   Matroid T1 = M1.remove(i), T2 = M2.remove(i);
   for (int j = 0; j < n; j++) if (!S[j]) {</pre>
    if (T1.can_add(j)) e[i].eb(j, -w[j]);
    if (T2.can_add(j)) e[j].eb(i, w[i]);
  } // maybe implicit build graph for more speed
  vector<pii> d(n + 2, \{INF, 0\}); d[n] = \{0, 0\};
  vector<int> prv(n + 2, -1);
  // change to SPFA for more speed, if necessary
  for (int upd = 1; upd--; )
   for (int u = 0; u < n + 2; u++)
    for (auto [v, c] : e[u]) {
     pii x(d[u].first + c, d[u].second + 1);
     if (x < d[v]) d[v] = x, prv[v] = u, upd = 1;
  if (d[n + 1].first >= INF) break;
  for (int x = prv[n+1]; x!=n; x = prv[x]) S.flip(x);
  // S is the max-weighted independent set w/ size sz
 return S;
    / from Nacl
```

Stable Marriage 8.3

```
1: Initialize m \in M and w \in W to free
2: while \exists free man m who has a woman w to propose to do
       w \leftarrow \text{first woman on } m's list to whom m has not yet proposed
       if \exists some pair (m', w) then
          if w prefers m to m' then
              m' \gets \mathit{free}
              (m, w) \leftarrow \mathsf{engaged}
8:
          end if
       else
```

```
(m, w) \leftarrow \mathsf{engaged}
11:
    end if
12: end while
8.4 Bitset LCS [4155ab]
cin >> n >> m;
for (int i = 1, x; i <= n; ++i)</pre>
cin >> x, p[x].set(i);
for (int i = 1, x; i <= m; ++i) {</pre>
cin >> x, (g = f) |= p[x];
f.shiftLeftByOne(), f.set(0);
 ((f = g - f)^{-} = g) \&= g;
cout << f.count() << '\n';</pre>
      Prefix Substring LCS [7d8faf]
void all_lcs(string S, string T) { // 0-base
vector<size_t> h(T.size()); iota(all(h), 1);
for (size_t a = 0; a < S.size(); ++a) {</pre>
 for (size_t c = 0, v = 0; c < T.size(); ++c)</pre>
   if (S[a] == T[c] || h[c] < v) swap(h[c], v);</pre>
  // here, LCS(s[0, a], t[b, c]) =
  // c - b + 1 - sum([h[i] > b] | i \le c)
} // test @ yosupo judge
     Convex 1D/1D DP [e5ab4b]
struct S { int i, l, r; };
auto solve(int n, int k, auto &w) {
vector < int64_t > dp(n + 1);
auto f = [&](int l, int r) -> int64_t {
 if (r - l > k) return -INF;
  return dp[l] + w(l + 1, r);
dp[0] = 0;
 deque<S> dq; dq.emplace_back(0, 1, n);
 for (int i = 1; i <= n; ++i) {
 dp[i] = f(dq.front().i, i);
 while (!dq.empty() && dq.front().r <= i)</pre>
  dq.pop_front();
  dq.front().l = i + 1;
 while (!dq.empty() &&
    f(i, dq.back().l) >= f(dq.back().i, dq.back().l))
   dq.pop_back();
  int p = i + 1;
  if (!dq.empty()) {
   auto [j, l, r] = dq.back();
   for (int s = 1 << 20; s; s >>= 1)
    if (l + s \le n \&\& f(i, l + s) \le f(j, l + s))
    l += s;
   dq.back().r = l; p = l + 1;
  if (p <= n) dq.emplace_back(i, p, n);</pre>
return dp;
} // test @ tioj 烏龜疊疊樂
     ConvexHull Optimization [b4318e]
mutable lld a, b, p;
bool operator<(const L &r) const {</pre>
  return a < r.a; /* here */ }
bool operator<(lld x) const { return p < x; }</pre>
lld Div(lld a, lld b) {
return a / b - ((a ^ b) < 0 && a % b); }
struct DynamicHull : multiset<L, less<>>> {
 static const lld kInf = 1e18;
bool Isect(iterator x, iterator y) {
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a)
  x->p = x->b > y->b ? kInf : -kInf; /* here */
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
void Insert(lld a, lld b) {
 auto z = insert({a, b, 0}), y = z++, x = y;
 while (Isect(y, z)) z = erase(z);
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
 while ((y = x) != begin() && (--x)->p >= y->p)
  Isect(x, erase(y));
lld Query(lld x) { // default chmax
  auto l = *lower_bound(x); // to chmin:
  return l.a * x + l.b; // modify the 2 "<>"
```

```
Min Plus Convolution [464dcd]
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector<int> min_plus_convolution(auto &a, auto &b) {
 const int n = (int)a.size(), m = (int)b.size();
 vector<int> c(n + m - 1, numeric_limits<int>::max());
 auto dc = [&](auto Y, int l, int r, int jl, int jr) {
  if (l > r) return;
  int mid = (l + r) / 2, from = -1, &best = c[mid];
  for (int j = jl; j <= jr; j++)
if (int i = mid - j; i >= 0 && i < n)</pre>
    if (best > a[i]+b[j]) best = a[i]+b[j], from = j;
  Y(Y, l, mid-1, jl, from); Y(Y, mid+1, r, from, jr);
 return dc(dc, 0, n-1+m-1, 0, m-1), c;
8.9 De-Bruijn [aa7700]
vector<int> de_bruijn(int k, int n) {
  // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
 vector<int> aux(n + 1), res;
 auto db = [&](auto self, int t, int p) -> void {
  if (t <= n)
   for (int i = aux[t - p]; i < k; ++i, p = t)</pre>
  aux[t] = i, self(self, t + 1, p);
else if (n % p == 0) for (int i = 1; i <= p; ++i)</pre>
   res.push_back(aux[i]);
 };
 return db(db, 1, 1), res;
8.10 Josephus Problem [7f9ceb]
lld f(lld n, lld m, lld k) { // n people kill m for
    each turn
 lld s = (m - 1) % (n - k); // O(k)
 for (lld i = n - k + 1; i <= n; i++) s = (s + m) % i;
 return s;
lld kth(lld n, lld m, i128 k) { // died at kth
 if (m == 1) return k; // O(m log(n))
 for (k = k*m+m-1; k >= n; k = k-n + (k-n)/(m-1));
 return k;
  // k and result are 0-based, test @ CF 101955
8.11 N Queens Problem [31f83e]
void solve(VI &ret, int n) { // no sol when n=2,3
 if (n % 6 == 2) {
  for (int i = 2; i <= n; i += 2) ret.push_back(i);</pre>
  ret.push_back(3); ret.push_back(1);
  for (int i = 7; i <= n; i += 2) ret.push_back(i);</pre>
  ret.push_back(5);
 } else if (n % 6 == 3) {
  for (int i = 4; i <= n; i += 2) ret.push_back(i);</pre>
  ret.push_back(2);
  for (int i = 5; i <= n; i += 2) ret.push_back(i);</pre>
  ret.push_back(1); ret.push_back(3);
 } else {
  for (int i = 2; i <= n; i += 2) ret.push_back(i);</pre>
  for (int i = 1; i <= n; i += 2) ret.push_back(i);</pre>
8.12 Tree Knapsack [f42766]
vector<int> G[N]; int dp[N][K]; pair<int,int> obj[N];
void dfs(int u, int mx) {
 for (int s : G[u]) {
  auto [w, v] = obj[s];
  if (mx < w) continue;</pre>
  for (int i = 0; i <= mx - w; i++)
   dp[s][i] = dp[u][i];
  dfs(s, mx - w);
  for (int i = w; i <= mx; i++)</pre>
   dp[u][i] = max(dp[u][i], dp[s][i - w] + v);
 }
8.13 Manhattan MST [1008bc]
vector<array<int, 3>> manhattanMST(vector<P> ps) {
 vector<int> id(ps.size()); iota(all(id), 0);
 vector<array<int, 3>> edges;
 for (int k = 0; k < 4; k++) {
  sort(all(id), [&](int i, int j) {
   return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y; });</pre>
  map<int, int> sweep;
```

```
for (int i : id) {
   for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++)) {
    if (P d = ps[i] - ps[it->second]; d.y > d.x) break;
   else edges.push_back({d.y + d.x, i, it->second});
  sweep[-ps[i].y] = i;
  for (P &p : ps)
   if (k \& 1) p.x = -p.x;
   else swap(p.x, p.y);
return edges; // [{w, i, j}, ...]
    test @ yosupo judge
8.14 Binary Search On Fraction [765c5a]
struct Q {
ll p, q;
Q go(Q b, ll d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
Q lo{0, 1}, hi{1, 0};
if (pred(lo)) return lo;
assert(pred(hi));
bool dir = 1, L = 1, H = 1;
for (; L || H; dir = !dir) {
 ll len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
  if (Q mid = hi.go(lo, len + step);
    mid.p > N || mid.q > N || dir ^ pred(mid))
    t++;
  else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
return dir ? hi : lo;
8.15 Barrett Reduction [d44617]
struct FastMod {
using Big = __uint128_t; llu b, m;
FastMod(llu b) : b(b), m(-1ULL / b) {}
llu reduce(llu a) { // a % b
 llu r = a - (llu)((Big(m) * a) >> 64) * b;
  return r >= b ? r - b : r;
8.16 Montgomery [648fb3]
struct Mont { // Montgomery multiplication
constexpr static int W = 64, L = 6;
llu mod, R1, R2, xinv;
void set_mod(llu _mod) {
 mod = _mod; assert(mod & 1); xinv = 1;
 for (int j = 0; j < L; j++) xinv *= 2 - xinv * mod;</pre>
 assert(xinv * mod == 1);
 const u128 R = (u128(1) << W) % mod;</pre>
 R1 = llu(R); R2 = llu(R*R \% mod);
llu redc(llu a, llu b) const {
 u128 T = u128(a) * b, m = -llu(T) * xinv;
 T += m * mod; T >>= W;
return llu(T >= mod ? T - mod : T);
llu from(llu x) const {
 assert(x < mod); return redc(x, R2); }</pre>
llu get(llu a) const { return redc(a, 1); }
llu one() const { return R1; }
} mont;
// a * b % mod == get(redc(from(a), from(b)))
```