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6 6 6 6 6 6 6 6 6 6 6 6 6	2 Segment & Line Intersection 3 2D Convex Hull 4 3D Convex Hull 5 2D Farthest Pair 6 kD Closest Pair (3D ver.) 7 Simulated Annealing 8 Half Plane Intersection 9 Minkowski Sum 10 Circle Class 11 Intersection of line and Circle 12 Intersection of Polygon and Circle	18 18 19 19 19 19 19 19 20 20	<pre>#define safe ((void)0) #define debug() ((void)0) #define orange() ((void)0) #endif  1.3 Increase Stack const int size = 256 &lt;&lt; 20; register long rsp asm("rsp"); char *p = (char*)malloc(size)+size, *bak = (char*)rsp;</pre>
6 6 6 6 6 6 6 6 6 6 6 6 6 6 6 6 6 6 6	2 Segment & Line Intersection 3 2D Convex Hull 4 3D Convex Hull 5 2D Farthest Pair 6 kD Closest Pair (3D ver.) 7 Simulated Annealing 8 Half Plane Intersection 9 Minkowski Sum 10 Circle Class 11 Intersection of line and Circle 12 Intersection of Polygon and Circle 13 Point & Hulls Tangent 14 Convex Hulls Tangent	18 18 19 19 19 19 19 19 20 20 20	<pre>#define safe ((void)0) #define debug() ((void)0) #define orange() ((void)0) #endif  1.3 Increase Stack  const int size = 256 &lt;&lt; 20; register long rsp asm("rsp"); char *p = (char*)malloc(size)+size, *bak = (char*)rsp;asm("movq %0, %%rsp\n"::"r"(p));</pre>
6 6 6 6 6 6 6 6 6 6 6 6	2 Segment & Line Intersection 3 2D Convex Hull 4 3D Convex Hull 5 2D Farthest Pair 6 kD Closest Pair (3D ver.) 7 Simulated Annealing 8 Half Plane Intersection 9 Minkowski Sum 10 Circle Class 11 Intersection of line and Circle 12 Intersection of Polygon and Circle 13 Point & Hulls Tangent 14 Convex Hulls Tangent 15 Tangent line of Two Circle	18 18 19 19 19 19 19 19 20 20 21 21	#define safe ((void)0) #define debug() ((void)0) #define orange() ((void)0) #endif  1.3 Increase Stack  const int size = 256 << 20; register long rsp asm("rsp"); char *p = (char*)malloc(size)+size, *bak = (char*)rsp;asm("movq %0, %%rsp\n"::"r"(p)); // main

## 1.4 Pragma Optimization

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse, sse2, sse3, ssse3, sse4")
#pragma GCC target("popent, abm, mmx, avx, tune=native")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8000)
```

## 1.5 IO Optimization

```
static inline int gc() {
constexpr int B = 1 << 20;
static char buf[B], *p, *q;
if(p == q \&\&
  (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
  return EOF;
 return *p++;
template < typename T >
static inline bool gn( T &x ) {
int c = gc(); T sgn = 1; x = 0;
while(('\theta'>c||c>'9') && c!=EOF && c!='-') c = gc(); if(c == '-') sgn = -1, c = gc();
if(c == EOF) return false;
while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
return x *= sgn, true;
```

#### 2 Data Structure

## 2.1 Dark Magic

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using namespace __gnu_pbds;
// heap tags: paring/binary/binomial/rc_binomial/thin
template<typename T>
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>, \
                  pairing_heap_tag>;
// pbds_heap::point_iterator
// x = pq.push(10); pq.modify(x, 87); a.join(b);
// tree tags: rb_tree_tag/ov_tree_tag/splay_tree_tag
template<typename T>
using ordered_set = tree<T, null_type, less<T>,
   rb_tree_tag, tree_order_statistics_node_update>;
// find_by_order, order_of_key
// hash tables: cc_hash_table/gp_hash_table
```

#### 2.2 Link-Cut Tree

```
template <typename Val, typename SVal> class LCT {
struct node
 int pa, ch[2];
 bool rev;
 Val v, prod, rprod;
 SVal sv, sub, vir
 node() : pa{0}, ch{0, 0}, rev{false}, v{}, prod{},
    rprod{}, sv{}, sub{}, vir{} {};
#define cur o[u]
#define lc cur.ch[0]
#define rc cur.ch[1]
vector<node> o;
bool is_root(int u) const {
 return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u;
bool is_rch(int u) const {
 return o[cur.pa].ch[1] == u && !is_root(u);
void down(int u) {
 if (not cur.rev) return;
 if (lc) set_rev(lc);
 if (rc) set_rev(rc);
 cur.rev = false;
void up(int u) {
 cur.prod = o[lc].prod * cur.v * o[rc].prod;
 cur.rprod = o[rc].rprod * cur.v * o[lc].rprod;
 cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
void set_rev(int u) {
 swap(lc, rc);
 swap(cur.prod, cur.rprod);
 cur.rev ^= 1;
```

```
void rotate(int u) {
 int f=cur.pa, g=o[f].pa, l=is_rch(u);
if (cur.ch[1 ^ 1]) o[cur.ch[1 ^ 1]].pa = f;
  if (not is_root(f)) o[g].ch[is_rch(f)] = u;
  o[f].ch[l] = cur.ch[l ^ 1];
  cur.ch[1 ^ 1] = f;
  cur.pa = g, o[f].pa = u;
 up(f);
 void splay(int u) {
  vector<int> stk = {u};
  while (not is_root(stk.back()))
   stk.push_back(o[stk.back()].pa);
  while (not stk.empty()) {
   down(stk.back());
   stk.pop_back();
  for (int f = cur.pa; not is_root(u); f = cur.pa) {
  if(!is_root(f))rotate(is_rch(u)==is_rch(f)?f:u);
   rotate(u);
 up(u);
 void access(int x) {
  for (int u = x, last = 0; u; u = cur.pa) {
   splay(u);
   cur.vir = cur.vir + o[rc].sub - o[last].sub;
   rc = last; up(last = u);
  }
  splay(x);
 int find_root(int u) {
  int la = 0;
  for (access(u); u; u = lc) down(la = u);
  return la;
 void split(int x, int y) {change_root(x);access(y);}
 void change_root(int u) { access(u); set_rev(u); }
public:
LCT(int n = 0) : o(n + 1) {}
 int add(const Val &v = {}) {
 o.push_back(v);
 return int(o.size()) - 2;
 int add(Val &&v) {
 o.emplace_back(move(v));
  return int(o.size()) - 2;
 void set_val(int u, const Val &v) {
  splay(++u); cur.v = v; up(u);
 void set_sval(int u, const SVal &v) {
 splay(++u); cur.sv = v; up(u);
 Val query(int x, int y) {
 split(++x, ++y); return o[y].prod;
 SVal subtree(int p, int u) {
 change_root(++p); access(++u);
  return cur.vir + cur.sv;
 bool connected(int u, int v) {
  return find_root(++u) == find_root(++v); }
 void link(int x, int y) {
 change_root(++x); access(++y);
 o[y].vir = o[y].vir + o[x].sub;
 up(o[x].pa = y);
 void cut(int x, int y) {
 split(++x, ++y);
o[y].ch[0] = o[x].pa = 0; up(y);
#undef cur
#undef lc
#undef rc
2.3 LiChao Segment Tree
```

```
struct L {
 int m, k, id;
 L(): id(-1) {}
L(int a, int b, int c) : m(a), k(b), id(c) {}
```

```
int at(int x) { return m * x + k; }
class LiChao {
private:
 int n; vector<L> nodes;
 static int lc(int x) { return 2 * x + 1; }
static int rc(int x) { return 2 * x + 2; }
 void insert(int 1, int r, int id, L ln) {
  int m = (1 + r) >> 1;
  if (nodes[id].id == -1) {
   nodes[id] = ln;
   return:
  bool atLeft = nodes[id].at(1) < ln.at(1);</pre>
  if (nodes[id].at(m) < ln.at(m)) {</pre>
   atLeft ^= 1
   swap(nodes[id], ln);
  if (r - 1 == 1) return;
  if (atLeft) insert(1, m, lc(id), ln);
  else insert(m, r, rc(id), ln);
 int query(int 1, int r, int id, int x) {
  int ret = 0, m = (1 + r) >> 1;
  if (nodes[id].id != -1)
   ret = nodes[id].at(x);
  if (r - 1 == 1) return ret;
  if (x < m) return max(ret, query(1, m, lc(id), x));</pre>
  return max(ret, query(m, r, rc(id), x));
public:
LiChao(int n_{-}) : n(n_{-}), nodes(n * 4) {}
 void insert(L ln) { insert(0, n, 0, ln); }
int query(int x) { return query(0, n, 0, x); }
```

## 2.4 Treap

```
namespace Treap{
#define sz(x)((x)?((x)-size):0)
struct node{
 int size;
 uint32_t pri;
 node *lc, *rc, *pa;
 \mathsf{node()} \colon\! \mathsf{size(0)}, \mathsf{pri}(\mathsf{rand())}, \mathsf{lc(0)}, \mathsf{rc(0)}, \mathsf{pa(0)} \{\}
 void pull() {
   size = 1; pa = nullptr;
  if ( lc ) { size += lc->size; lc->pa = this; }
   if ( rc ) { size += rc->size; rc->pa = this; }
}:
node* merge( node* L, node* R ) {
 if ( not L or not R ) return L ? L : R;
 if ( L->pri > R->pri ) {
  L->rc = merge( L->rc, R ); L->pull();
  return L:
  } else {
   R->lc = merge( L, R->lc ); R->pull();
   return R;
void split_by_size( node*rt,int k,node*&L,node*&R ) {
 if ( not rt ) L = R = nullptr;
 else if( sz( rt->lc ) + 1 <= k ) {
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
  .
L->pull();
  } else {
  R = rt;
   split_by_size( rt->lc, k, L, R->lc );
   R->pull();
 } // sz(L) == k
int getRank(node *o) { // 1-base
 int r = sz(o->lc) + 1;
 for (;o->pa != nullptr; o = o->pa)
   if (o->pa->rc == o) r += sz(o->pa->lc) + 1;
  return r;
#undef sz
```

## 2.5 Linear Basis

```
template <int BITS, typename S = int> struct Basis {
 static constexpr auto MIN = numeric_limits<S>::min();
 array<pair<uint64_t, S>, BITS> b;
 Basis() { b.fill({0, MIN}); }
 void add(uint64_t x, S p) {
  for (int i = BITS-1; i>=0; i--) if ((x >> i) & 1) {
    if (b[i].first == 0) return b[i]={x, p}, void();
    if (b[i].second < p)</pre>
     swap(b[i].first, x), swap(b[i].second, p);
    x ^= b[i].first;
  }
 optional<uint64_t> query_kth(uint64_t v, uint64_t k){
  vector<pair<uint64_t, int>> o;
  for (int i = 0; i < BITS; i++)</pre>
    if (b[i].first) o.emplace_back(b[i].first, i);
  if (k >= (1ULL << o.size())) return {};
for (int i = int(o.size()) - 1; i >= 0; i--)
    if ((k >> i & 1) ^ (v >> o[i].second & 1))
     v ^= o[i].first;
  return v;
 Basis filter(S 1) {
  Basis res = *this;
for (int i = 0; i < BITS; i++)</pre>
    if (res.b[i].second < 1) res.b[i] = {0, MIN};</pre>
  return res:
};
```

## 2.6 Binary Search On Segment Tree

```
// find_first = x \rightarrow minimal x s.t. check([a, x))
// find_last = x \rightarrow maximal x s.t. check([x, b))
template <typename C>
int find_first(int 1, const C &check) {
 if (1 >= n) return n + 1;
 1 += sz;
 for (int i = height; i > 0; i--)
  propagate(1 >> i);
 Monoid sum = identity;
 do {
  while ((1 & 1) == 0) 1 >>= 1;
  if (check(f(sum, data[1]))) {
   while (1 < sz) {
    propagate(1);
    1 <<= 1;
    auto nxt = f(sum, data[1]);
    if (not check(nxt)) {
     sum = nxt;
     1++:
    }
   return 1 + 1 - sz;
  sum = f(sum, data[1++]);
 } while ((1 & -1) != 1);
 return n + 1;
template <typename C>
int find_last(int r, const C &check) {
 if (r <= 0) return -1;</pre>
 r += sz;
 for (int i = height; i > 0; i--)
  propagate((r - 1) >> i);
 Monoid sum = identity;
 do {
  while (r > 1 and (r & 1)) r >>= 1;
  if (check(f(data[r], sum))) {
   while (r < sz) {</pre>
    propagate(r);
    r = (r << 1) + 1;
    auto nxt = f(data[r], sum);
    if (not check(nxt)) {
     sum = nxt:
     r--;
    }
   return r - sz;
```

} else {

```
National Taiwan University - ckiseki
  sum = f(data[r], sum);
                                                                ++ch, dfs(v, u);
                                                                low[u] = min(low[u], low[v]);
 } while ((r & -r) != r);
                                                                if (low[v] > dfn[u])
return -1;
                                                                bridge[t] = true
                                                               if (low[v] >= dfn[u])
                                                                ap[u] = true;
3
    Graph
    2-SAT (SCC)
                                                              ap[u] &= (ch != 1 or u != f);
class TwoSat{
private:
                                                           public:
int n;
                                                             void init(int n_) {
vector<vector<int>> rG,G,sccs;
                                                              g.assign(n = n_, vector<pair<int, int>>());
vector<int> ord,idx;
                                                              low.assign(n, ecnt = 0);
vector<bool> vis,result;
                                                             dfn.assign(n, 0);
void dfs(int u){
                                                             ap.assign(n, false);
 vis[u]=true
 for(int v:G[u])
                                                             void add_edge(int u, int v) {
  if(!vis[v]) dfs(v);
                                                             g[u].emplace_back(v, ecnt);
 ord.push_back(u);
                                                             g[v].emplace_back(u, ecnt++);
void rdfs(int u){
                                                             void solve() {
 vis[u]=false;idx[u]=sccs.size()-1;
                                                             bridge.assign(ecnt, false);
  sccs.back().push_back(u);
                                                             for (int i = 0; i < n; ++i)
 for(int v:rG[u])
                                                               if (not dfn[i]) dfs(i, i);
  if(vis[v])rdfs(v);
                                                             bool is_ap(int x) { return ap[x]; }
public:
                                                            bool is_bridge(int x) { return bridge[x]; }
 void init(int n_){
 G.clear();G.resize(n=n_);
  rG.clear();rG.resize(n);
                                                           3.3 Round Square Tree
  sccs.clear();ord.clear();
                                                           int N, M, cnt;
 idx.resize(n);result.resize(n);
                                                           std::vector<int> G[maxn], T[maxn * 2];
void add_edge(int u,int v){
                                                           int dfn[maxn], low[maxn], dfc;
 G[u].push_back(v);rG[v].push_back(u);
                                                           int stk[maxn], tp;
void orr(int x,int y){
                                                            void Tarjan(int u) {
 if ((x^y)==1)return
                                                            low[u] = dfn[u] = ++dfc;
 add_edge(x^1,y); add_edge(y^1,x);
                                                             stk[++tp] = u;
                                                             for (int v : G[u]) {
bool solve(){
                                                             if (!dfn[v]) {
 vis.clear();vis.resize(n);
for(int i=0;i<n;++i)</pre>
                                                              Tarjan(v)
                                                               low[u] = std::min(low[u], low[v]);
   if(not vis[i])dfs(i);
                                                               if (low[v] == dfn[u]) {
  reverse(ord.begin(),ord.end());
                                                                ++cnt
  for (int u:ord){
                                                               for (int x = 0; x != v; --tp) {
  if(!vis[u])continue;
                                                                x = stk[tp];
   sccs.push_back(vector<int>());
                                                                T[cnt].push_back(x);
   rdfs(u);
                                                                T[x].push_back(cnt);
  for(int i=0;i<n;i+=2)</pre>
                                                                T[cnt].push_back(u);
   if(idx[i]==idx[i+1])
                                                               T[u].push_back(cnt);
   return false:
  vector<bool> c(sccs.size());
                                                              } else
  for(size_t i=0;i<sccs.size();++i){</pre>
                                                               low[u] = std::min(low[u], dfn[v]);
  for(auto sij : sccs[i]){
    result[sij]=c[i]
    c[idx[sij^1]]=!c[i];
                                                           int main() { // ...
                                                            cnt = N;
 return true;
                                                            for (int u = 1; u <= N; ++u)</pre>
                                                             if (!dfn[u]) Tarjan(u), --tp;
bool get(int x){return result[x];}
int get_id(int x){return idx[x];}
int count(){return sccs.size();}
                                                           3.4 Centroid Decomposition
} sat2;
                                                           struct Centroid {
                                                             vector<vector<int64_t>> Dist;
3.2 BCC
                                                             vector<int> Pa, Dep;
class BCC {
                                                             vector<int64_t> Sub, Sub2;
private:
                                                             vector<int> Cnt, Cnt2;
                                                             vector<int> vis, sz, mx, tmp
int n, ecnt;
vector<vector<pair<int, int>>> g;
                                                             void DfsSz(int x) +
vector<int> dfn, low;
                                                              vis[x] = true; sz[x] = 1; mx[x] = 0;
                                                              for (auto [u, w] : g[x]) {
vector<bool> ap, bridge;
void dfs(int u, int f) {
                                                               if (vis[u]) continue;
 dfn[u] = low[u] = dfn[f] + 1;
                                                              DfsSz(u);
  int ch = 0;
                                                               sz[x] += sz[u];
  for (auto [v, t] : g[u]) if (v != f) {
                                                              mx[x] = max(mx[x], sz[u]);
   if (dfn[v]) {
   low[u] = min(low[u], dfn[v]);
                                                              tmp.push_back(x);
```

// find cycle

ans += in[i];

**if** (vis[x] == i) ·

id[y] = tot; id[x] = tot++;

if (!tot) return ans;

for (auto &e : E) {

for (int i = 0; i < n; i++)
if (id[i] == -1) id[i] = tot++;</pre>

vector<int> id(n, -1), vis(n, -1); for (int i = 0; i < n; i++) {</pre>

for (int x = i; x != -1 && id[x] == -1; x = prv[x])

for (int y = prv[x]; y != x; y = prv[y])

int tot = 0

break;

}

vis[x] = i;

```
void DfsDist(int x, int64_t D = 0) {
                                                                     if (id[e.u] != id[e.v]) e.w -= in[e.v];
 Dist[x].push_back(D); vis[x] = true;
for (auto [u, w] : g[x])
                                                                     e.u = id[e.u], e.v = id[e.v];
   if (not vis[u]) DfsDist(u, D + w);
                                                                    n = tot; root = id[root];
                                                                   }
                                                                 }
 void DfsCen(int x, int D = 0, int p = -1) {
  tmp.clear(); DfsSz(x);
                                                                } DMST;
  int M = tmp.size();
                                                                3.6 Dominator Tree
  int C = -1;
  for (int u : tmp) {
  if (max(M - sz[u], mx[u]) * 2 <= M) C = u;</pre>
   vis[u] = false;
  DfsDist(C);
                                                                void init(int n) {
  for (int u : tmp) vis[u] = false;
 Pa[C] = p; vis[C] = true; Dep[C] = D;
for (auto [u, w] : g[C])
   if (not vis[u]) DfsCen(u, D + 1, C);
 Centroid(int N, vector<vector<pair<int,int>>> g)
                                                                 for (int i = 0; i < n; ++i) {</pre>
  : Sub(N), Sub2(N), Cnt(N), Cnt2(N), Dist(N),
  Pa(N), Dep(N), vis(N), sz(N), mx(N)
  { DfsCen(0); }
 void Mark(int v) {
  int x = v, z = -1;
                                                                void dfs(int x) {
                                                                 rev[dfn[x] = tk] = x;
  for (int i = Dep[v]; i >= 0; --i) {
   Sub[x] += Dist[v][i]; Cnt[x]++;
   if (z != -1) {
    Sub2[z] += Dist[v][i];
    Cnt2[z]++;
                                                                   r[dfn[u]].push_back(dfn[x]);
   z = x; x = Pa[x];
                                                                }
                                                                int find(int x, int c = 0) {
 int64_t Query(int v) {
                                                                  int p = find(fa[x], 1);
  int64_t res = 0;
  int x = v, z = -1;
 for (int i = Dep[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Cnt[x] * Dist[v][i];
                                                                 fa[x] = p;
return c ? p : val[x];
  if (z != -1) res-=Sub2[z]+1LL*Cnt2[z]*Dist[v][i];
   z = x; x = Pa[x];
  return res;
                                                                  dfs(s);
3.5 Directed Minimum Spanning Tree
struct Edge { int u, v, w; };
struct DirectedMST { // find maximum
                                                                   for (int &u : rdom[i]) {
                                                                    int p = find(u);
 int solve(vector<Edge> E, int root, int n) {
                                                                    if (sdom[p] == i) dom[u] = i;
else dom[u] = p;
  int ans = 0;
  while (true) {
   // find best in edge
   vector<int> in(n, -inf), prv(n, -1);
                                                                   if (i) merge(i, rp[i]);
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
     in[e.v] = e.w;
                                                                  for (int i = 1; i < tk; ++i)</pre>
     prv[e.v] = e.u;
   in[root] = 0; prv[root] = -1;
   for (int i = 0; i < n; i++)</pre>
    if (in[i] == -inf) return -inf;
```

# namespace dominator { vector<int> g[maxn], r[maxn], rdom[maxn]; int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn]; int dom[maxn], val[maxn], rp[maxn], tk; // vertices are numbered from $\theta$ to n-1fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);fill(fa, fa + n, -1); fill(val, val + n, -1);fill(sdom, sdom + n, -1); fill(rp, rp + n, -1); fill(dom, dom + n, -1); tk = 0; g[i].clear(); r[i].clear(); rdom[i].clear(); void add\_edge(int x, int y) { g[x].push\_back(y); } fa[tk] = sdom[tk] = val[tk] = tk; tk ++;for (int u : g[x]) { if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x]; void merge(int x, int y) { fa[x] = y; } if (fa[x] == x) return c ? -1 : x; if (p == -1) return c ? fa[x] : val[x]; if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]]; vector<int> build(int s, int n) { // return the father of each node in the dominator tree // p[i] = -2 if i is unreachable from sfor (int i = tk - 1; i >= 0; --i) { for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]); if (i) rdom[sdom[i]].push\_back(i); vector<int> p(n, -2); p[s] = -1; if (sdom[i] != dom[i]) dom[i] = dom[dom[i]]; for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre> 3.7 Edge Coloring $// \max(d_u) + 1$ edge coloring, time: O(NM)int C[kN][kN], G[kN][kN]; // 1-based, G: ans void clear(int N) { for (int i = 0; i <= N; i++) for (int j = 0; j <= N; j++) C[i][j] = G[i][j] = 0;</pre> void solve(vector<pair<int, int>> &E, int N) { int X[kN] = {}, a; auto update = [&](int u) for (X[u] = 1; C[u][X[u]]; X[u]++); auto color = [&](int u, int v, int c) { **int** p = G[u][v] G[u][v] = G[v][u] = c;C[u][c] = v, C[v][c] = u;C[u][p] = C[v][p] = 0;if (p) X[u] = X[v] = pelse update(u), update(v);

```
return p;
auto flip = [&](int u, int c1, int c2) {
 int p = C[u][c1];
 swap(C[u][c1], C[u][c2]);
 if (p) G[u][p] = G[p][u] = c2;
 if (!C[u][c1]) X[u] = c1;
 if (!C[u][c2]) X[u] = c2;
for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {</pre>
 auto [u, v] = E[t];
 int v0 = v, c = X[u], c0 = c, d;
 vector<pair<int, int>> L; int vst[kN] = {};
 while (!G[u][v0]) {
  L.emplace_back(v, d = X[v]);
  if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
    c = color(u, L[a].first, c);
  else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
    color(u, L[a].first, L[a].second);
  else if (vst[d]) break
  else vst[d] = 1, v = C[u][d];
 if (!G[u][v0]) {
  for (; v; v = flip(v, c, d), swap(c, d));
if (C[u][c0]) { a = int(L.size()) - 1;
   while (--a >= 0 && L[a].second != c);
   for(;a>=0;a--)color(u,L[a].first,L[a].second);
  } else t--;
```

#### **Lowbit Decomposition** 3.8

```
class LBD {
int timer, chains;
vector<vector<int>> G;
vector<int> tl, tr, chain, head, dep, pa;
// chains : number of chain
// tl, tr[u] : subtree interval in the seq. of u
// head[i] : head of the chain i
// chian[u] : chain id of the chain u is on
void predfs(int u, int f) {
 dep[u] = dep[pa[u] = f] + 1;
  for (int v : G[u]) if (v != f) {
  predfs(v, u);
   if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
    chain[u] = chain[v];
  if (chain[u] == 0) chain[u] = ++chains;
void dfschain(int u, int f) {
 tl[u] = timer++;
 if (head[chain[u]] == -1)
  head[chain[u]] = u;
  for (int v : G[u])
  if (v != f and chain[v] == chain[u])
    dfschain(v, u)
  for (int v : G[u])
   if (v != f and chain[v] != chain[u])
    dfschain(v, u);
  tr[u] = timer;
public:
LBD(int n) : timer(0), chains(0), G(n), tl(n), tr(n),
chain(n), head(n, -1), dep(n), pa(n) {}
void_add_edge(int u, int v) {
 G[u].push_back(v); G[v].push_back(u);
void decompose() { predfs(0, 0); dfschain(0, 0); }
PII get_subtree(int u) { return {tl[u], tr[u]}; }
vector<PII> get_path(int u, int v) {
 vector<PII> res;
  while (chain[u] != chain[v]) {
  if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
    swap(u, v)
   int s = head[chain[u]];
   res.emplace_back(tl[s], tl[u] + 1);
   u = pa[s];
  if (dep[u] < dep[v]) swap(u, v);</pre>
```

```
return res:
 }
};
```

res.emplace\_back(tl[v], tl[u] + 1);

## Manhattan Minimum Spanning Tree

```
typedef Point<int> P;
vector<array<int, 3>> manhattanMST(vector<P> ps) {
 vi id(sz(ps));
 iota(all(id), 0);
 vector<array<int, 3>> edges;
 rep(k, 0, 4) {
  sort(all(id), [&](int i, int j) {
   return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
  }):
  map<int, int> sweep;
  for (int i : id) {
   for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++)) {
    int j = it->second;
    P d = ps[i] - ps[j];
    if (d.y > d.x) break;
    edges.push_back({d.y + d.x, i, j});
   sweep[-ps[i].y] = i;
  for (P &p : ps)
   if (k \& 1) p.x = -p.x;
   else swap(p.x, p.y);
 return edges; // [{w, i, j}, ...]
}
```

#### 3.10 MaxClique

```
// contain a self loop u to u, than u won't in clique
template < size_t MAXN >
class MaxClique{
private:
 using bits = bitset< MAXN >;
 bits popped, G[ MAXN ], ans
 size_t deg[ MAXN ], deo[ MAXN ], n;
 void sort_by_degree() {
  popped.reset();
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
    deg[ i ] = G[ i ].count();
  for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;
    for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )</pre>
        mi = deg[id = j];
    popped[ deo[ i ] = id ] = 1;
    for( size_t u = G[ i ]._Find_first() ;
     u < n ; u = G[ i ]._Find_next( u ) )
      -- deg[ u ];
 void BK( bits R, bits P, bits X ) {
  if (R.count()+P.count() <= ans.count()) return;</pre>
  if ( not P.count() and not X.count() )
   if ( R.count() > ans.count() ) ans = R;
   return:
  /* greedily chosse max degree as pivot
  bits cur = P \mid X; size_t pivot = \theta, sz = \theta;
  for ( size_t u = cur._Find_first() ;
   u < n ; u = cur._Find_next( u )</pre>
    if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
  cur = P & ( ~G[ pivot ] );
  */ // or simply choose first
  bits cur = P & (~G[ ( P | X )._Find_first() ]);
  for ( size_t u = cur._Find_first()
          ; u = cur._Find_next( u ) ) {
   if ( R[ u ] ) continue;
   R[u] = 1;
   BK( R, P & G[ u ], X & G[ u ] );
   R[u] = P[u] = 0, X[u] = 1;
  }
public:
 void init( size_t n_ ) {
  for ( size_t i = 0 ; i < n ; ++ i )
```

vector<int> r, c;

```
G[ i ].reset();
                                                                  for (int i = 0; i < n; i++)
                                                                   if (mask[i]) r.push_back(i);
  ans.reset();
                                                                  for (int i = 0; i < n; i++)</pre>
 void add_edges( int u, bits S ) { G[ u ] = S; }
void add_edge( int u, int v ) {
                                                                   d[i] = int((a[i] & mask).count());
                                                                  sort(r.begin(), r.end(),
  G[u][v] = G[v][u] = 1;
                                                                   [&](int i, int j) { return d[i] > d[j]; });
                                                                  csort(r, c);
 int solve() {
                                                                  dfs(r, c, 1, mask);
  sort_by_degree(); // or simply iota( deo... )
                                                                  return ans; // sol[0 ~ ans-1]
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
   deg[ i ] = G[ i ].count();
                                                               } graph;
  bits pob, nob = 0; pob.set();
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
                                                               3.12 Minimum Mean Cycle
  for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
                                                               /* minimum mean cycle O(VE) */
   size_t v = deo[ i ];
                                                               struct MMC{
   bits tmp; tmp[ v ] = 1;
                                                               #define FZ(n) memset((n),0,sizeof(n))
  BK( tmp, pob & G[ v ], nob & G[ v ] );
pob[ v ] = 0, nob[ v ] = 1;
                                                               #define E 101010
                                                               #define V 1021
                                                               #define inf 1e9
  return static_cast< int >( ans.count() );
                                                                 struct Edge { int v,u; double c; };
                                                                 int n, m,
                                                                           prv[V][V], prve[V][V], vst[V];
};
                                                                 Edge e[E];
                                                                 vector<int> edgeID, cycle, rho;
      MaxCliqueDyn
3.11
                                                                 double d[V][V];
constexpr int kN = 150;
                                                                 void init( int _n ) { n = _n; m = 0; }
struct MaxClique { // Maximum Clique
                                                                 // WARNING: TYPE matters
 bitset<kN> a[kN], cs[kN]
                                                                 void add_edge( int vi , int ui , double ci )
{ e[ m ++ ] = { vi , ui , ci }; }
 int ans, sol[kN], q, cur[kN], d[kN], n;
 void init(int _n) {
                                                                 void bellman_ford() {
                                                                 for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);
  for(int i=0; i++);</pre>
 n = _n, ans = q = 0;
  for (int i = 0; i < n; i++) a[i].reset();</pre>
 void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
                                                                   for(int j=0; j<m; j++)</pre>
 void csort(vector<int> &r, vector<int> &c) {
                                                                    int v = e[j].v, u = e[j].u;
  int mx = 1, km = max(ans - q + 1, 1), t = 0,
                                                                    if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
    m = int(r.size());
                                                                     d[i+1][u] = d[i][v]+e[j].c;
  cs[1].reset(); cs[2].reset();
                                                                     prv[i+1][u] = v;
  for (int i = 0; i < m; i++) {
                                                                     prve[i+1][u] = j;
   int p = r[i], k = 1;
while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
                                                                  }
   cs[k][p] = 1;
   if (k < km) r[t++] = p;
                                                                 double solve(){
                                                                  // returns inf if no cycle, mmc otherwise
  c.resize(m);
                                                                  double mmc=inf;
  if(t) c[t-1] = 0;
                                                                  int st = -1;
  for (int k = km; k <= mx; k++) {</pre>
                                                                  bellman_ford();
   for (int p = int(cs[k]._Find_first());
                                                                  for(int i=0; i<n; i++) {</pre>
      p < kN; p = int(cs[k]._Find_next(p))) {</pre>
                                                                   double avg=-inf;
    r[t] = p; c[t++] = k;
                                                                   for(int k=0; k<n; k++) {</pre>
   }
                                                                    if(d[n][i]<inf-eps)</pre>
  }
                                                                     avg=max(avg,(d[n][i]-d[k][i])/(n-k));
 }
                                                                    else avg=max(avg,inf);
 void dfs(vector<int> &r, vector<int> &c, int 1,
  bitset<kN> mask) {
                                                                   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
  while (!r.empty()) {
   int p = r.back(); r.pop_back();
                                                                  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
   mask[p] = 0;
                                                                  for (int i=n; !vst[st]; st=prv[i--][st]) {
   if (q + c.back() <= ans) return;</pre>
                                                                   vst[st]++
   cur[q++] = p;
                                                                   edgeID.PB(prve[i][st]);
   vector<int> nr, nc;
                                                                   rho.PB(st);
   bitset<kN> nmask = mask & a[p];
   for (int i : r)
                                                                  while (vst[st] != 2) {
    if (a[p][i]) nr.push_back(i);
                                                                  int v = rho.back(); rho.pop_back();
   if (!nr.empty()) {
                                                                   cycle.PB(v);
    if (1 < 4) {
                                                                   vst[v]++;
     for (int i : nr)
      d[i] = int((a[i] & nmask).count());
                                                                  reverse(ALL(edgeID));
     sort(nr.begin(), nr.end(),
                                                                  edgeID.resize(SZ(cycle));
      [&](int x, int y)
                                                                  return mmc:
       return d[x] > d[y];
      });
                                                               } mmc;
   csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
                                                                       Mo's Algorithm on Tree
    ans = q; copy(cur, cur + q, sol);
                                                               dfs u:
                                                                 push u
   c.pop_back(); q--;
                                                                 iterate subtree
  }
                                                                 push u
                                                               Let P = LCA(u, v) with St(u) \le St(v)
 int solve(bitset<kN> mask) { // vertex mask
                                                               if (P == u) query[St(u), St(v)]
```

else query[Ed(u), St(v)], query[St(P), St(P)]

#### 3.14 Virtual Tree

```
vector<pair<int, int>> build(vector<int> vs, int r) {
vector<pair<int, int>> res;
sort(vs.begin(), vs.end(), [](int i, int j) {
  return dfn[i] < dfn[j]; });</pre>
 vector<int> s = {r};
for (int v : vs) if (v != r) {
 if (int o = lca(v, s.back()); o != s.back()) {
   while (s.size() >= 2) {
    if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
    res.emplace_back(s[s.size() - 2], s.back());
    s.pop_back();
  if (s.back() != o) {
    res.emplace_back(o, s.back());
    s.back() = o;
  s.push_back(v);
for (size_t i = 1; i < s.size(); ++i)</pre>
 res.emplace_back(s[i - 1], s[i]);
 return res; // (x, y): x->y
```

## 4 Matching & Flow

## 4.1 Bipartite Matching

```
struct BipartiteMatching {
vector<int> X[N];
int fX[N], fY[N], n;
bitset<N> vis;
bool dfs(int x)
  for (auto i : X[x]) if (not vis[i]) {
  vis[i] = true;
  if (fY[i] == -1 || dfs(fY[i])) {
    fY[fX[x] = i] = x;
    return true:
  return false;
void init(int n_, int m) {
 fill_n(X, n = n_, vector<int>());
 memset(fX, -1, sizeof(int) * n);
 memset(fY, -1, sizeof(int) * m);
void add_edge(int x, int y) { X[x].push_back(y); }
int solve() { // return how many pair matched
 int cnt = 0;
 for (int i = 0; i < n; i++) {
  vis.reset();
  cnt += dfs(i);
  return cnt;
};
```

## 4.2 Dijkstra Cost Flow

```
// kN = \#(vertices)
// MCMF.{Init, AddEdge, MincostMaxflow}
// MincostMaxflow(source, sink, flow_limit, &cost)
// => flow
using Pii = pair<int, int>;
constexpr int kInf = 0x3f3f3f3f, kN = 500;
struct Edge {
int to, rev, cost, flow;
struct MCMF { // 0-based
int n{}, m{}, s{}, t{};
vector<Edge> graph[kN];
 // Larger range for relabeling
int64_t dis[kN] = {}, h[kN] = {};
int p[kN] = {};
void Init(int nn) {
 n = nn;
 for (int i = 0; i < n; i++) graph[i].clear();</pre>
void AddEdge(int u, int v, int f, int c) {
 graph[u].push_back({v,
   static_cast<int>(graph[v].size()), c, f});
```

```
graph[v].push_back(
   {u, static_cast<int>(graph[u].size()) - 1,
    -c, 0});
 bool Dijkstra(int &max_flow, int64_t &cost) {
  priority_queue<Pii, vector<Pii>, greater<>> pq;
  fill_n(dis, n, kInf);
  dis[s] = 0;
  pq.emplace(0, s);
  while (!pq.empty()) {
   auto u = pq.top();
   pq.pop();
   int v = u.second;
   if (dis[v] < u.first) continue;</pre>
   for (auto &e : graph[v]) {
    auto new_dis =
     dis[v] + e.cost + h[v] - h[e.to];
    if (e.flow > 0 && dis[e.to] > new_dis) {
     dis[e.to] = new_dis;
     p[e.to] = e.rev;
     pq.emplace(dis[e.to], e.to);
  if (dis[t] == kInf) return false;
  for (int i = 0; i < n; i++) h[i] += dis[i];</pre>
  int d = max_flow;
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   d = min(d, graph[e.to][e.rev].flow);
  max_flow -= d;
  cost += int64_t(d) * h[t];
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   e.flow += d;
   graph[e.to][e.rev].flow -= d;
  return true;
 int MincostMaxflow(
  int ss, int tt, int max_flow, int64_t &cost) {
  this->s = ss, this->t = tt;
  cost = 0;
  fill_n(h, n, 0);
  auto orig_max_flow = max_flow;
  while (Dijkstra(max_flow, cost) && max_flow) {}
  return orig_max_flow - max_flow;
};
4.3 Dinic
```

```
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
    int to, rev;
    Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
      int u = bfs.front(); bfs.pop();
      for (auto e: G[u]) {
        if (e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
        bfs.push(e.to); lv[e.to] = lv[u] + 1;
      }
    return lv[ed] != -1;
  Cap DFS(int u, Cap f){
    if (u == ed) return f;
    Cap ret = 0;
    for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
      auto &e = G[u][i];
```

```
if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;
        Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
        G[e.to][e.rev].cap += nf;
        if (f == 0) return ret;
     if (ret == 0) lv[u] = -1;
     return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
  G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
  st = st_, ed = ed_; Cap ret = 0;
     while (BFS()) {
        idx.assign(n, 0);
        Cap f = DFS(st, numeric_limits<Cap>::max());
        ret += f;
        if (f == 0) break;
     return ret;
  }
};
```

### 4.4 Construct VC

```
vi cover(vector<vi>& g, int n, int m) {
  vi match(m, -1);
  int res = dfsMatching(g, match);
  vector<bool> lfound(n, true), seen(m);
for (int it : match) if (it != -1) lfound[it] = false
  vi q, cover;
  rep(i,0,n) if (lfound[i]) q.push_back(i);
  while (!q.empty()) {
    int i = q.back(); q.pop_back();
    lfound[i] = 1;
    for (int e : g[i]) if (!seen[e] && match[e] != -1)
       seen[e] = true;
       q.push_back(match[e]);
  rep(i,0,n) if (!lfound[i]) cover.push_back(i);
rep(i,0,m) if (seen[i]) cover.push_back(n+i);
  assert(sz(cover) == res);
  return cover;
```

## Flow Models

- · Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.

  - 2. For each edge (x,y,l,u), connect  $x\to y$  with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower
  - 4. If in(v) > 0, connect  $S \to v$  with capacity in(v), otherwise, con- $\text{nect } v \to T \text{ with capacity } -in(v).$ 
    - To maximize, connect t 
      ightarrow s with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Oth-
    - erwise, the maximum flow from s to t is the answer To minimize, let f be the maximum flow from S to T. Connect  $t\,\rightarrow\,s$  with capacity  $\infty$  and let the flow from S to Tbe f'. If  $f+f'\neq \sum_{v\in V, in(v)>0}in(v)$  , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e + f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
  - 1. Redirect every edge:  $y \to x$  if  $(x, y) \in M$ ,  $x \to y$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source S and sink T
  - 2. For each edge (x, y, c), connect  $x \to y$  with (cost, cap) = (c, 1) if c>0, otherwise connect y o x with (cost, cap)=(-c,1)
  - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v)>0, connect S o v with (cost,cap)=0(0, d(v))

- 5. For each vertex v with d(v) < 0, connect v  $\rightarrow$  T with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let K be the sum of all weights
  - 3. Connect source  $s \to v$ ,  $v \in G$  with capacity K
  - 4. For each edge (u, v, w) in G, connect  $u \to v$  and  $v \to u$  with
  - 5. For  $v \in {\it G}$ , connect it with sink  $v \to t$  with capacity K+2T-1 $\left(\sum_{e \in E(v)} w(e)\right) - 2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \rightarrow v'$  with
  - weight w(u,v). 2. Connect v o v' with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - Find the minimum weight perfect matching on G'.
- Project selection problem
  - 1. If  $p_v\,>\,0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create
  - edge (v,t) with capacity  $-p_v$ . 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
  - 3. The mincut is equivalent to the maximum profit of a subset of

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity  $c_x$  and create edge (s,y) with capacity  $c_y$ .
- 2. Create edge (x,y) with capacity  $c_{xy}$ . 3. Create edge (x,y) and edge (x',y') with capacity  $c_{xyx'y'}$ .

## 4.6 General Graph Matching namespace matching {

```
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
 for (int i = 0; i \le n; ++i) match[i] = pre[i] = n;
 for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
 g[u].push_back(v);
 g[v].push_back(u);
int Find(int u) {
return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
 static int tk = 0; tk++;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
void Blossom(int x, int y, int 1) {
while (Find(x) != 1) {
  pre[x] = y, y = match[x];
  if (s[y] == 1) q.push(y), s[y] = 0;
  if (fa[x] == x) fa[x] = 1;
if (fa[y] == y) fa[y] = 1;
  x = pre[y];
 }
bool Bfs(int r, int n) {
 for (int i = 0; i \le n; ++i) fa[i] = i, s[i] = -1;
 while (!q.empty()) q.pop();
 q.push(r);
 s[r] = 0;
 while (!q.empty()) {
  int x = q.front(); q.pop();
  for (int u : g[x]) {
   if (s[u] == -1) {
pre[u] = x, s[u] = 1;
    if (match[u] == n) {
```

```
for (int a = u, b = x, last; b != n; a = last, b =
    pre[a])
    last = match[b], match[b] = a, match[a] = b;
    return true;
}
    q.push(match[u]);
    s[match[u]] = 0;
} else if (!s[u] && Find(u) != Find(x)) {
    int 1 = LCA(u, x, n);
    Blossom(x, u, 1);
    Blossom(u, x, 1);
}
}
return false;
}
int Solve(int n) {
    int res = 0;
    for (int x = 0; x < n; ++x) {
        if (match[x] == n) res += Bfs(x, n);
    }
    return res;
}</pre>
```

## 4.7 Global Min-Cut

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
pair<int, int> phase(int n) {
memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
int s = -1, t = -1;
 while (true) {
  int c = -1;
  for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
   if (c == -1 \mid | g[i] > g[c]) c = i;
  if (c == -1) break;
  v[s = t, t = c] = true;
  for (int i = 0; i < n; ++i) {</pre>
   if (del[i] || v[i]) continue;
   g[i] += w[c][i];
  }
 return make_pair(s, t);
int mincut(int n) {
int cut = 1e9:
 memset(del, false, sizeof(del));
 for (int i = 0; i < n - 1; ++i) {
 int s, t; tie(s, t) = phase(n);
  del[t] = true; cut = min(cut, g[t]);
  for (int j = 0; j < n; ++j) {
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
 return cut;
```

### 4.8 GomoryHu Tree

```
int g[maxn];
vector<edge> GomoryHu(int n) {
  vector<edge> rt;
  for(int i=1;i<=n;++i)g[i]=1;
  for(int i=2;i<=n;++i) {
    int t=g[i];
    flow.reset(); // clear flows on all edge
    rt.push_back({i,t,flow(i,t)});
    flow.walk(i); // bfs points that connected to i (use
        edges not fully flow)
    for(int j=i+1;j<=n;++j) {
        if(g[j]==t && flow.connect(j))g[j]=i; // check if i
        can reach j
    }
  }
  return rt;
}</pre>
```

## 4.9 Kuhn Munkres

```
class KM {
private:
 static constexpr lld INF = 1LL << 60;</pre>
 vector<lld> hl,hr,slk;
 vector<int> fl,fr,pre,qu;
 vector<vector<lld>> w;
 vector<bool> v1,vr;
 int n, ql, qr;
 bool check(int x) {
  if (v1[x] = true, f1[x] != -1)
   return vr[qu[qr++] = f1[x]] = true;
  while (x != -1) swap(x, fr[fl[x] = pre[x]]);
  return false:
 void bfs(int s) {
  fill(slk.begin(), slk.end(), INF);
  fill(v1.begin(), v1.end(), false);
fill(vr.begin(), vr.end(), false);
  ql = qr = 0;
  vr[qu[qr++] = s] = true;
  while (true) {
   11d d;
   while (ql < qr) {</pre>
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
   if(!v1[x]&&s1k[x]>=(d=h1[x]+hr[y]-w[x][y])){
      if (pre[x] = y, d) slk[x] = d;
       else if (!check(x)) return;
    }
    }
   d = INF;
   for (int x = 0; x < n; ++x)
if (!v1[x] && d > s1k[x]) d = s1k[x];
    for (int x = 0; x < n; ++x) {
    if (v1[x]) h1[x] += d;
     else slk[x] -= d;
    if (vr[x]) hr[x] -= d;
   for (int x = 0; x < n; ++x)
    if (!v1[x] && !slk[x] && !check(x)) return;
public:
 void init( int n_ ) {
  qu.resize(n = n_);
  fl.assign(n, -1); fr.assign(n, -1);
hr.assign(n, 0); hl.resize(n);
  w.assign(n, vector<lld>(n));
  slk.resize(n); pre.resize(n);
  vl.resize(n); vr.resize(n);
 void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 11d solve() {
  for (int i = 0; i < n; ++i)
   hl[i] = *max_element(w[i].begin(), w[i].end());
  for (int i = 0; i < n; ++i) bfs(i);</pre>
  11d res = 0;
  for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
  return res;
} km;
```

### 4.10 Minimum Cost Circulation

```
while(!mark[upd])mark[upd]=1,upd=pv[upd];
     return upd:
    }
    idx++:
 }
return -1;
int Solve(int n) {
int rt = -1, ans = 0;
while ((rt = NegativeCycle(n)) >= 0) {
 memset(mark, false, sizeof(mark));
 vector<pair<int, int>> cyc;
 while (!mark[rt]) {
  cyc.emplace_back(pv[rt], ed[rt]);
  mark[rt] = true;
  rt = pv[rt];
 reverse(cyc.begin(), cyc.end());
 int cap = kInf;
 for (auto &i : cyc)
  auto &e = g[i.first][i.second];
  cap = min(cap, e.cap);
 for (auto &i : cyc)
  auto &e = g[i.first][i.second];
  e.cap -= cap;
  g[e.to][e.rev].cap += cap;
  ans += e.cost * cap;
 }
return ans;
```

## 4.11 Minimum Cost Maximum Flow

```
class MiniCostMaxiFlow{
using Cap = int; using Wei = int64_t;
using PCW = pair<Cap,Wei>;
static constexpr Cap INF_CAP = 1 << 30;</pre>
static constexpr Wei INF_WEI = 1LL<<60;</pre>
private:
struct Edge{
 int to, back;
 Cap cap; Wei wei;
 Edge() {}
 Edge(int a,int b, Cap c, Wei d):
   to(a),back(b),cap(c),wei(d) {}
};
int ori, edd;
vector<vector<Edge>> G;
vector<int> fa, wh;
vector<bool> inq;
vector<Wei> dis;
PCW SPFA(){
 fill(inq.begin(),inq.end(),false);
 fill(dis.begin(), dis.end(), INF_WEI);
  queue<int> qq; qq.push(ori);
 dis[ori] = 0;
 while(not qq.empty()){
  int u=qq.front();qq.pop();
  ing[u] = false;
   for(int i=0;i<SZ(G[u]);++i){</pre>
   Edge e=G[u][i];
    int v=e.to; Wei d=e.wei;
    if(e.cap <= 0 | |dis[v] <= dis[u] + d)
     continue
    dis[v] = dis[u] + d;
   fa[v] = u, wh[v] = i;
   if (ing[v]) continue;
    qq.push(v);
    inq[v] = true;
  }
  if(dis[edd]==INF_WEI) return {-1, -1};
 Cap mw=INF_CAP;
 for(int i=edd;i!=ori;i=fa[i])
  mw=min(mw,G[fa[i]][wh[i]].cap);
 for (int i=edd;i!=ori;i=fa[i]){
  auto &eg=G[fa[i]][wh[i]];
  eg.cap -= mw;
```

```
G[eg.to][eg.back].cap+=mw;
  return {mw, dis[edd]};
public:
 void init(int n){
  G.clear();G.resize(n);
  fa.resize(n);wh.resize(n);
  inq.resize(n); dis.resize(n);
 void add_edge(int st, int ed, Cap c, Wei w){
  G[st].emplace_back(ed,SZ(G[ed]),c,w);
  G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
 PCW solve(int a, int b){
  ori = a, edd = b;
Cap cc=0; Wei ww=0;
  while(true){
   PCW ret=SPFA();
   if(ret.first==-1) break;
   cc+=ret.first:
   ww+=ret.first * ret.second;
  return {cc,ww};
} mcmf;
```

```
4.12 Maximum Weight Graph Matching
struct WeightGraph {
 static const int inf = INT_MAX;
 static const int maxn = 514;
 struct edge {
  int u, v, w;
 edge(){}
  edge(int u, int v, int w): u(u), v(v), w(w) {}
 int n, n_x;
 edge g[maxn * 2][maxn * 2];
 int lab[maxn * 2];
 int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa
    [maxn * 2];
 int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
    maxn * 2]
 vector<int> flo[maxn * 2];
 queue<int> q;
 int e_delta(const edge &e) { return lab[e.u] + lab[e.v
    ] - g[e.u][e.v].w * 2; }
 void update_slack(int u, int x) { if (!slack[x] ||
    e_delta(g[u][x]) < e_delta(g[slack[x]][x])) slack[x
    ] = u; }
 void set_slack(int x) {
 slack[x] = 0;
  for (int u = 1; u <= n; ++u)</pre>
   if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
    update_slack(u, x);
 void q_push(int x) {
 if (x \le n) q.push(x);
  else for (size_t i = 0; i < flo[x].size(); i++)</pre>
    q_push(flo[x][i]);
 void set_st(int x, int b) {
 st[x] = b;
  if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)
     set_st(flo[x][i], b);
 int get_pr(int b, int xr) {
 int pr = find(flo[b].begin(), flo[b].end(), xr) - flo
    [b].begin();
  if (pr % 2 == 1)
  reverse(flo[b].begin() + 1, flo[b].end());
   return (int)flo[b].size() - pr;
  return pr;
 void set_match(int u, int v) {
 match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u, xr)
  for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo</pre>
    [u][i ^ 1]);
```

```
set_match(xr, v);
 rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
                                                               }
   end());
void augment(int u, int v) {
 for (; ; ) {
  int xnv = st[match[u]];
  set_match(u, v);
  if (!xnv) return;
  set_match(xnv, st[pa[xnv]]);
  u = st[pa[xnv]], v = xnv;
 }
int get_lca(int u, int v) {
static int t = 0;
 for (++t; u || v; swap(u, v)) {
  if (u == 0) continue;
  if (vis[u] == t) return u;
  vis[u] = t;
  u = st[match[u]];
  if (u) u = st[pa[u]];
 }
 return 0;
void add_blossom(int u, int lca, int v) {
 int b = n + 1;
 while (b <= n_x && st[b]) ++b;</pre>
 if (b > n_x) ++n_x;
 lab[b] = 0, S[b] = 0
 match[b] = match[lca];
 flo[b].clear();
 flo[b].push_back(lca);
 for (int x = u, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
 reverse(flo[b].begin() + 1, flo[b].end())
 for (int x = v, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
 set_st(b, b);
 for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].w
 for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;</pre>
 for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  for (int x = 1; x <= n_x; ++x)
   if (g[b][x].w == 0 \mid \mid e_delta(g[xs][x]) < e_delta(g[xs][x])
   [b][x]))
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
  for (int x = 1; x <= n; ++x)
   if (flo_from[xs][x]) flo_from[b][x] = xs;
 set_slack(b);
void expand_blossom(int b) {
 for (size_t i = 0; i < flo[b].size(); ++i)</pre>
  set_st(flo[b][i], flo[b][i]);
 int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
   xr):
 for (int i = 0; i < pr; i += 2)
  int xs = flo[b][i], xns = flo[b][i + 1];
  pa[xs] = g[xns][xs].u;
  S[xs] = 1, S[xns] = 0;
  slack[xs] = 0, set_slack(xns);
  q_push(xns);
 S[xr] = 1, pa[xr] = pa[b];
 for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  S[xs] = -1, set_slack(xs);
 st[b] = 0;
bool on_found_edge(const edge &e) {
 int u = st[e.u], v = st[e.v];
 if (S[v] == -1) {
  pa[v] = e.u, S[v] = 1
  int nu = st[match[v]]
  slack[v] = slack[nu] = 0;
 S[nu] = 0, q_push(nu);
} else if (S[v] == 0) {
  int lca = get_lca(u, v);
```

```
if (!lca) return augment(u,v), augment(v,u), true;
  else add_blossom(u, lca, v);
 return false;
bool matching() {
 \begin{array}{lll} memset(S+1, -1, & sizeof(int) * n_x); \\ memset(slack + 1, 0, sizeof(int) * n_x); \end{array} 
 q = queue<int>();
for (int x = 1; x <= n_x; ++x)
if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0,
   q_push(x);
 if (q.empty()) return false;
  while (q.size()) {
   int u = q.front(); q.pop();
   if (S[st[u]] == 1) continue
   for (int v = 1; v <= n; ++v)
    if (g[u][v].w > 0 && st[u] != st[v]) {
     if (e_delta(g[u][v]) == 0) {
      if (on_found_edge(g[u][v])) return true;
     } else update_slack(u, st[v]);
  int d = inf;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b && S[b] == 1) d = min(d, lab[b] / 2)
  for (int x = 1; x <= n_x; ++x)
if (st[x] == x && slack[x]) {</pre>
    if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x])
    else if (S[x] == 0) d = min(d, e_delta(g[slack[x
   ]][x]) / 2);
  for (int u = 1; u <= n; ++u) {
   if (S[st[u]] == 0) {
    if (lab[u] <= d) return 0;</pre>
    lab[u] -= d;
   } else if (S[st[u]] == 1) lab[u] += d;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b) {
    if (S[st[b]] == 0) lab[b] += d * 2;
    else if (S[st[b]] == 1) lab[b] -= d * 2;
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
if (st[x] == x && slack[x] && st[slack[x]] != x &&</pre>
   e_{delta}(g[slack[x]][x]) == 0)
    if (on_found_edge(g[slack[x]][x])) return true;
  for (int b = n + 1; b <= n_x; ++b)
if (st[b] == b && S[b] == 1 && lab[b] == 0)
   expand_blossom(b);
 return false;
pair<long long, int> solve() {
memset(match + 1, 0, sizeof(int) * n);
 n_x = n:
 int n_matches = 0;
 long long tot_weight = 0;
 for (int u = 0; u \le n; ++u) st[u] = u, flo[u].clear
   ();
 int w_max = 0;
 for (int u = 1; u <= n; ++u)</pre>
 for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);</pre>
   w_max = max(w_max, g[u][v].w);
 for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
 while (matching()) ++n_matches;
 for (int u = 1; u <= n; ++u)
  if (match[u] && match[u] < u)</pre>
   tot_weight += g[u][match[u]].w;
 return make_pair(tot_weight, n_matches);
void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
   g[vi][ui].w = wi; }
void init(int _n) {
 n = _n;
for (int u = 1; u <= n; ++u)
```

```
for (int v = 1; v <= n; ++v)
g[u][v] = edge(u, v, 0);
```

#### 5 Math

#### 5.1 **Common Bounds**

### 5.1.1 Partition function

$$\begin{split} p(0) &= 1, \; p(n) = \sum_{k \in \mathbb{Z} \backslash \{0\}} (-1)^{k+1} p(n-k(3k-1)/2) \\ & \qquad \qquad p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n}) \\ & \qquad \qquad \frac{n}{p(n)} \; \frac{0\,12\,3\,4\,5\,6\,7\,8\,9\,20\,50\,100}{11\,12\,3\,5\,7\,11\,15\,22\,30\,627\,{\sim}2e5\,{\sim}2e8} \end{split}$$

## 5.1.2 Divisor function

n	100	1e3	1e6	1e9	1e12	1e15	1e18
$\max_{i \leq n} (d(i))$	12	32	240	1344	6720	26880	103680

## 5.1.3 Factorial

n	123	4 5	6	7	8	9	9	10	
n!	126	24 12	0 720	5040	3 403	20 362	880 3	62880	0
n	11								
n!						.3e12 2			
n						100			, ,
n!	2e18	2e25	3e32	8e47	3e64	9e157	6e262	2 > DB	L_MAX

## 5.1.4 Binom Coef

## Strling Number

#### 5.2.1 First Kind

 $S_1(n,k)$  counts the number of permutations of n elements with k disjoint cvcles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1)\dots(x+n-1) = \sum_{k=0}^n S_1(n,k)x^k$$

$$g(x) = x(x+1)\dots(x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)!a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

#### 5.2.2 Second Kind

 $S_2(n,k)$  counts the number of ways to partition a set of n elements into knonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

## ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 \ (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
if (y == 0) g=x,a=1,b=0;
else exgcd(y,x%y,g,b,a),b=(x/y)*a;
```

#### 5.4 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
 vector<T> d(output.size() + 1), me, he;
 for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
 for (size_t j = 0; j < me.size(); ++j)
d[i] += output[i - j - 2] * me[j];</pre>
  if ((d[i] -= output[i - 1]) == 0) continue;
  if (me.empty()) {
   me.resize(f = i);
   continue;
  vector<T> o(i - f - 1);
  T k = -d[i] / d[f]; o.push_back(-k);
  for (T x : he) o.push_back(x * k);
  if (o.size() < me.size()) o.resize(me.size());</pre>
  for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
  if (i-f+he.size() >= me.size()) he = me, f = i;
  me = o;
 return me;
```

## 5.5 Charateristic Polynomial

```
vector<vector<int>> Hessenberg(const vector<vector<int
    >> &A) {
 int N = A.size();
 vector<vector<int>> H = A;
 for (int i = 0; i < N - 2; ++i) {
  if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {
    if (H[j][i]) {
     for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
     for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k
    ][j]);
     break:
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
   * H[i + 1][k] * (kP - coef)) % kP;
for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
 return H;
vector<int> CharacteristicPoly(const vector<vector<int
    >> &A) {
 int N = A.size();
 auto H = Hessenberg(A);
 for (int i = 0; i < N; ++i) {
 for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
 vector<vector<int>> P(N + 1, vector<int>(N + 1));
 P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
  P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
    1];
  int val = 1;
  for (int j = i - 1; j >= 0; --j) {
  int coef = 1LL * val * H[j][i - 1] % kP;
   for (int k = 0; k <= j; ++k) P[i][k] = (P[i][k] + 1
LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 if (N & 1) {
 for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
 return P[N];
```

#### 5.6 Chinese Remainder

```
x = a1 % m1
x = a2 % m2
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
// 0 <= x < lcm(m1, m2)</pre>
```

## 5.7 De-Bruijn

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
  if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
} else {
 aux[t] = aux[t - p];
  db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {</pre>
  aux[t] = i;
   db(t + 1, t, n, k);
}
int de_bruijn(int k, int n) {
  // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
if (k == 1) {
 res[0] = 0;
  return 1;
for (int i = 0; i < k * n; i++) aux[i] = 0;</pre>
db(1, 1, n, k);
 return sz;
```

## 5.8 DiscreteLog

```
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
  // x^? \equiv y (mod M)
  Int t = 1, c = 0, g = 1;
  for (Int M_ = M; M_ > 0; M_ >>= 1)
   g = g * x % M;
  for (g = gcd(g, M); t % g != 0; ++c) {
   if (t == y) return c;
t = t * x % M;
  if (y % g != 0) return -1;
 t /= g, y /= g, M /= g;
 Int h = 0, gs = 1;

for (; h * h < M; ++h) gs = gs * x % M;
 unordered_map<Int, Int> bs;
 for (Int s = 0; s < h; bs[y] = ++s)
   y = y * x % M;
  for (Int s = 0; s < M; s += h) {
   t = t * gs % M;
    if (bs.count(t)) return c + s + h - bs[t];
  return -1;
```

### 5.9 Extended Euler

```
a^b \equiv \begin{cases} a^{(b \mod \varphi(m)) + \varphi(m)} & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m
```

### 5.10 ExtendedFloorSum

```
\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}
```

```
\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}
```

```
5.11 Fast Fourier Transform
const int mod = 1000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
  static_assert (M1 <= M2 && M2 <= M3);
  constexpr int64_t r12 = modpow(M1, M2-2, M2);
  constexpr int64_t r13 = modpow(M1, M3-2, M3);
  constexpr int64_t r23 = modpow(M2, M3-2, M3);
  constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
  B = (B - A + M2) * r12 % M2;
  C = (C - A + M3) * r13 % M3;
  C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
 for (int i = 0; i <= maxn; i++)</pre>
  omega[i] = cplx(cos(2 * pi * j / maxn),
     sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
 int z = __builtin_ctz(n) - 1;
for (int i = 0; i < n; ++i) {</pre>
  int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^{=(i >> j & 1) << (z - j);
  if (x > i) swap(v[x], v[i]);
 for (int s = 2; s <= n; s <<= 1) {
  int z = s >> 1;
for (int i = 0; i < n; i += s) {</pre>
   for (int k = 0; k < z; ++k) {
    cplx x = v[i + z + k] * omega[maxn / s * k];
    v[i + z + k] = v[i + k] - x;
    v[i+k] = v[i+k] + x;
  }
void ifft(vector<cplx> &v, int n) {
 fft(v, n); reverse(v.begin() + 1, v.end());
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
VL convolution(const VI &a, const VI &b) {
 // Should be able to handle N <= 10^5, C <= 10^4
 int sz = 1;
 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 vector<cplx> v(sz);
 for (int i = 0; i < sz; ++i) {
  double re = i < a.size() ? a[i] : 0;
double im = i < b.size() ? b[i] : 0;</pre>
  v[i] = cplx(re, im);
 fft(v, sz);
 for (int i = 0; i <= sz / 2; ++i) {
  int j = (sz - i) & (sz - 1);
  cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
    * cplx(0, -0.25);
  if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
    ].conj()) * cplx(0, -0.25);
  v[i] = x;
```

ifft(v, sz);

VL c(sz);

```
for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
                                                                            x1), inv = (x0-x1, x1) w/o final div */
                                                              * x = (x0+x1,
                                                            void fwt(int x[], int N, bool inv = false) {
return c:
                                                               for (int d = 1; d < N; d <<= 1) {
VI convolution_mod(const VI &a, const VI &b, int p) {
                                                                 for (int s = 0, d2 = d * 2; s < N; s += d2)
                                                                   for (int i = s, j = s + d; i < s + d; i++, j++) {
int sz = 1:
while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
                                                                     int ta = x[i], tb = x[j];
 vector<cplx> fa(sz), fb(sz);
                                                                     x[i] = modadd(ta, tb);
for (int i = 0; i < (int)a.size(); ++i)</pre>
                                                                     x[j] = modsub(ta, tb);
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 for (int i = 0; i < (int)b.size(); ++i)</pre>
 fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
                                                               if (inv) for (int i = 0, invn = modinv(N); i < N; i</pre>
 fft(fa, sz), fft(fb, sz);
double r = 0.25 / sz;
                                                                 x[i] = modmul(x[i], invn);
 cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
 int j = (sz - i) & (sz - 1);
cplx a1 = (fa[i] + fa[j].conj());
                                                            5.14 Miller Rabin
                                                            bool isprime(llu x) {
  static auto witn = [](llu a, llu u, llu n, int t) {
 cplx a2 = (fa[i] - fa[j].conj()) * r2;
 cplx b1 = (fb[i] + fb[j].conj()) * r3;
                                                               if (!a) return false;
  cplx b2 = (fb[i] - fb[j].conj()) * r4;
                                                               while (t--) {
  if (i != j) {
                                                               llu a2 = mmul(a, a, n);
  cplx c1 = (fa[j] + fa[i].conj());
                                                                if (a2 == 1 && a != 1 && a != n - 1) return true;
   cplx c2 = (fa[j] - fa[i].conj()) * r2;
                                                                a = a2;
   cplx d1 = (fb[j] + fb[i].conj()) * r3;
                                                               }
   cplx d2 = (fb[j] - fb[i].conj()) * r4;
                                                               return a != 1;
   fa[i] = c1 * d1 + c2 * d2 * r5;
   fb[i] = c1 * d2 + c2 * d1;
                                                              if (x < 2) return false;</pre>
                                                              if (!(x & 1)) return x == 2;
  fa[j] = a1 * b1 + a2 * b2 * r5;
                                                              int t = __builtin_ctzll(x - 1);
  fb[j] = a1 * b2 + a2 * b1;
                                                              llu \ odd = (x - 1) >> t;
                                                             for (llu m:
  {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
fft(fa, sz), fft(fb, sz);
vector<int> res(sz);
                                                               if (witn(mpow(m % x, odd, x), odd, x, t))
for (int i = 0; i < sz; ++i) {</pre>
                                                               return false:
 long long a = round(fa[i].re), b = round(fb[i].re),
                                                              return true;
       c = round(fa[i].im);
 res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
}
                                                            5.15 NTT
 return res;
                                                            template <int mod, int G, int maxn>
}}
                                                             struct NTT {
5.12 FloorSum
                                                              static_assert (maxn == (maxn & -maxn));
                                                              int roots[maxn];
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
                                                              NTT () {
                                                               int r = modpow(G, (mod - 1) / maxn);
// @return sum_\{i=\theta\}^{n-1} floor((ai + b)/m) mod 2^64
                                                               for (int i = maxn >> 1; i; i >>= 1) {
1lu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
                                                                roots[i] = 1;
11u ans = 0;
                                                                for (int j = 1; j < i; j++)</pre>
while (true) {
                                                                 roots[i + j] = modmul(roots[i + j - 1], r);
 if (a >= m) {
                                                                r = modmul(r, r);
  ans += n * (n - 1) / 2 * (a / m); a %= m;
  if (b >= m) {
                                                              // n must be 2^k, and 0 <= F[i] < mod
  ans += n * (b / m); b %= m;
                                                              void operator()(int F[], int n, bool inv = false) {
                                                               for (int i = 0, j = 0; i < n; i++) {
 llu y_max = a * n + b;
                                                               if (i < j) swap(F[i], F[j]);</pre>
 if (y_max < m) break;</pre>
                                                                for (int k = n>1; (j^k < k; k>=1);
 // y_max < m * (n + 1)
 // floor(y_max / m) <= n
                                                               for (int s = 1; s < n; s *= 2) {
 n = (11u)(y_max / m), b = (11u)(y_max % m);
                                                                for (int i = 0; i < n; i += s * 2) {
 swap(m, a);
                                                                 for (int j = 0; j < s; j++) {
                                                                  int a = F[i+j]
return ans;
                                                                  int b = modmul(F[i+j+s], roots[s+j]);
                                                                  F[i+j] = modadd(a, b); // a + b
11d floor_sum(lld n, lld m, lld a, lld b) {
                                                                  F[i+j+s] = modsub(a, b); // a - b
llu ans = 0;
if (a < 0) {
 11u a2 = (a \% m + m) \% m;
 ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
                                                               if (inv) {
 a = a2:
                                                               int invn = modinv(n);
                                                                for (int i = 0; i < n; i++)</pre>
if (b < 0) {
                                                                F[i] = modmul(F[i], invn);
 11u b2 = (b \% m + m) \% m;
                                                                reverse(F + 1, F + n);
 ans -= 1ULL * n * ((b2 - b) / m);
 b = b2;
return ans + floor_sum_unsigned(n, m, a, b);
                                                            NTT<2013265921, 31, 1048576> ntt;
5.13 FWT
                                                            5.16 Partition Number
/* or convolution:
                                                            int b = sqrt(n)
* x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                            ans[0] = tmp[0] = 1;
* and convolution:
                                                            for (int i = 1; i <= b; i++) {
```

#define fi(1, r) for (int i = int(1); i < int(r); ++i)

```
for (int rep = 0; rep < 2; rep++)</pre>
                                                              template <int mod, int G, int maxn> struct Poly : V {
  for (int j = i; j <= n - i * i; j++)
                                                               static uint32_t n2k(uint32_t n) {
   modadd(tmp[j], tmp[j-i]);
                                                                if (n <= 1) return 1;
 for (int j = i * i; j <= n; j++)
modadd(ans[j], tmp[j - i * i]);</pre>
                                                                return 1u << (32 - __builtin_clz(n - 1));</pre>
                                                               static NTT<mod, G, maxn> ntt; // coefficients in [0, P)
                                                               explicit Poly(int n = 1) : V(n) {}
      Pi Count (Linear Sieve)
                                                               Poly(const V &v) : V(v) {}
                                                               Poly(const Poly &p, size_t n) : V(n) {
static constexpr int N = 1000000 + 5;
                                                                copy_n(p.data(), min(p.size(), n), data());
11d pi[N];
vector<int> primes;
                                                               Poly &irev() { return reverse(data(), data() + size())
bool sieved[N];
                                                                    *this; }
                                                               Poly &isz(int sz) { return resize(sz), *this; }
11d cube_root(11d x){
 lld s=cbrt(x-static_cast<long double>(0.1));
                                                               Poly &iadd(const Poly &rhs) { // n() == rhs.n()
 while(s*s*s <= x) ++s;</pre>
                                                                fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
 return s-1;
                                                                return *this;
11d square_root(11d x){
                                                               Poly &imul(int k) {
 1ld s=sqrt(x-static_cast<long double>(0.1));
                                                                fi(0, size())(*this)[i] = modmul((*this)[i], k);
 while(s*s <= x) ++s;</pre>
                                                                return *this:
 return s-1;
                                                               Poly Mul(const Poly &rhs) const {
void init(){
                                                                const int sz = n2k(size() + rhs.size() - 1);
primes.reserve(N)
                                                                Poly X(*this, sz), Y(rhs, sz);
ntt(X.data(), sz), ntt(Y.data()
 primes.push_back(1);
                                                                                                   sz):
 for(int i=2;i<N;i++) {</pre>
                                                                fi(0, sz) X[i] = modmul(X[i], Y[i]);
  if(!sieved[i]) primes.push_back(i);
                                                                ntt(X.data(), sz, true);
  pi[i] = !sieved[i] + pi[i-1];
                                                                return X.isz(size() + rhs.size() - 1);
  for(int p: primes) if(p > 1) {
  if(p * i >= N) break;
                                                               Poly Inv() const { // coef[0] != 0
   sieved[p * i] = true;
                                                                if (size() == 1) return V{modinv(*begin())};
   if(p % i == 0) break;
                                                                const int sz = n2k(size() * 2);
                                                                Poly X = Poly(*this, (size() + 1) / 2).Inv().isz(sz),
 }
                                                                    Y(*this, sz);
                                                                ntt(X.data(), sz), ntt(Y.data(), sz);
11d phi(11d m, 11d n) {
 static constexpr int MM = 80000, NN = 500;
                                                                fi(0, sz) X[i] = modmul(X[i], modsub(2, modmul(X[i],
                                                                  Y[i])));
 static lld val[MM][NN];
                                                                ntt(X.data(), sz, true);
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
                                                                return X.isz(size());
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
                                                               Poly Sqrt() const { // coef[0] \in [1, mod)^2
 lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
                                                                if (size() == 1) return V{QuadraticResidue((*this)
 if(m<MM&&n<NN) val[m][n] = ret+1;
                                                                  [0], mod)};
 return ret;
                                                                Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
                                                                  size());
11d pi_count(11d);
                                                                return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
11d P2(11d m, 11d n) {
                                                                  + 1):
 11d sm = square_root(m), ret = 0;
 for(lld i = n+1;primes[i]<=sm;i++)</pre>
                                                               pair<Poly, Poly> DivMod(const Poly &rhs) const {
  ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
                                                                if (size() < rhs.size()) return {V{0}, *this};</pre>
 return ret;
                                                                const int sz = size() - rhs.size() + 1;
                                                                Poly X(rhs); X.irev().isz(sz);
11d pi_count(11d m) {
                                                                Poly Y(*this); Y.irev().isz(sz);
 if(m < N) return pi[m];</pre>
                                                                Poly Q = Y.Mul(X.Inv()).isz(sz).irev();
 11d n = pi_count(cube_root(m));
                                                                X = rhs.Mul(Q), Y = *this;
 return phi(m, n) + n - 1 - P2(m, n);
                                                                fi(0, size()) Y[i] = modsub(Y[i], X[i]);
                                                                return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
       Pollard Rho
                                                               Poly Dx() const {
// does not work when n is prime
                                                                Poly ret(size() - 1);
// return any non-trivial factor
                                                                fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
llu pollard_rho(llu n) {
 static auto f = [](llu x, llu k, llu m) {
                                                                return ret.isz(max<int>(1, ret.size()));
    return add(k, mul(x, x, m), m); };
 if (!(n & 1)) return 2;
                                                               Poly Sx() const {
 mt19937 rnd(120821011);
                                                                Poly ret(size() + 1);
 while (true) {
                                                                fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
 llu y = 2, yy = y, x = rnd() % n, t = 1;
for (llu sz = 2; t == 1; sz <<= 1, y = yy) {
                                                                  this)[i]);
                                                                return ret;
   for (llu i = 0; t == 1 && i < sz; ++i) {
    yy = f(yy, x, n);
                                                               Poly Ln() const { // coef[0] == 1
    t = gcd(yy > y ? yy - y : y - yy, n);
                                                                return Dx().Mul(Inv()).Sx().isz(size());
   }
                                                               Poly Exp() const \{ // coef[\theta] == \theta \}
  if (t != 1 && t != n) return t;
                                                                if (size() == 1) return V{1};
                                                                Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size)
                                                                  ()):
                                                                Poly Y = X.Ln(); Y[0] = mod - 1;
      Polynomial Operations
5.19
                                                                fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
return X.Mul(Y).isz(size());
using V = vector<int>;
```

// maximize c^Tx under Ax <= B

```
Poly Pow(const string &K) const {
                                                                  // return VD(n, -inf) if the solution doesn't exist
                                                                  // return VD(n, +inf) if the solution is unbounded
  int nz = 0;
  while (nz < size() && !(*this)[nz]) ++nz;</pre>
                                                                  using VD = vector<double>;
  int nk = 0, nk2 = 0;
                                                                  using VVD = vector<vector<double>>;
  for (char c : K) {
                                                                  const double eps = 1e-9;
   nk = (nk * 10 + c - '0') % mod;
                                                                  const double inf = 1e+9;
   nk2 = nk2 * 10 + c - '0';
                                                                   int n, m;
   if (nk2 * nz >= size())
                                                                  VVD d:
                                                                  vector<int> p, q;
void pivot(int r, int s) {
    return Poly(size());
   nk2 %= mod - 1;
                                                                    double inv = 1.0 / d[r][s];
  if (!nk && !nk2) return Poly(V{1}, size());
                                                                    for (int i = 0; i < m + 2; ++i)
                                                                    for (int j = 0; j < n + 2; ++j)
if (i != r && j != s)
  Poly X = V(data() + nz, data() + size() - nz * (nk2 - nz)
     1));
  int x0 = X[0];
                                                                       d[i][j] -= d[r][j] * d[i][s] * inv;
                                                                    for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
  return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
    modpow(x0, nk2)).irev().isz(size()).irev();
                                                                    d[r][s] = inv; swap(p[r], q[s]);
 Poly InvMod(int L) { // (to evaluate linear recursion)
  Poly R{1, 0}; // *this * R mod x^L = 1 (*this[0] ==
                                                                  bool phase(int z) {
                                                                   int x = m + z
  for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                                    while (true) {
   Poly 0 = R.Mul(Poly(data(), min<int>(2 << level,
                                                                     int s = -1;
    size())))
                                                                     for (int i = 0; i <= n; ++i) {</pre>
   Poly Q(2 << level); Q[0] = 1;
                                                                      if (!z && q[i] == -1) continue;
   for (int j = (1 << level); j < (2 << level); ++j)</pre>
                                                                      if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
    Q[j] = modsub(mod, O[j]);
   R = R.Mul(Q).isz(4 << level);</pre>
                                                                     if (d[x][s] > -eps) return true;
                                                                     int r = -1;
for (int i = 0; i < m; ++i) {</pre>
  }
  return R.isz(L);
                                                                      if (d[i][s] < eps) continue;</pre>
 static int LinearRecursion(const V &a, const V &c,
                                                                      if (r == -1 ||
    int64_t n) { // a_n = \sum c_j a_(n-j)}
                                                                       d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  const int k = (int)a.size();
                                                                     if (r == -1) return false;
  assert((int)c.size() == k + 1);
  Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
                                                                    pivot(r, s);
  fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
  C[k] = 1
  while (n) {
                                                                  VD solve(const VVD &a, const VD &b, const VD &c) {
                                                                   m = b.size(), n = c.size();
   if (n % 2) W = W.Mul(M).DivMod(C).second;
   n /= 2, M = M.Mul(M).DivMod(C).second;
                                                                    d = VVD(m + 2, VD(n + 2))
                                                                    for (int i = 0; i < m; ++i)</pre>
                                                                    for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
  int ret = 0;
  fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
                                                                    p.resize(m), q.resize(n + 1);
                                                                    for (int i = 0; i < m; ++i) 
p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
  return ret:
                                                                    for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
                                                                    q[n] = -1, d[m + 1][n] = 1;
#undef fi
using Poly_t = Poly<998244353, 3, 1 << 20>;
                                                                    int r = 0;
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
                                                                    for (int i = 1; i < m; ++i)</pre>
                                                                     if (d[i][n + 1] < d[r][n + 1]) r = i;
5.20 Quadratic residue
                                                                    if (d[r][n + 1] < -eps) {</pre>
struct S {
                                                                     pivot(r, n);
 int MOD, w;
                                                                     if (!phase(1) || d[m + 1][n + 1] < -eps)
 int64_t x, y;
                                                                      return VD(n, -inf);
                                                                     for (int i = 0; i < m; ++i) if (p[i] == -1) {
 S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
                                                                      int s = min_element(d[i].begin(), d[i].end() - 1)
  : MOD(m), w(w_{-}), x(x_{-}), y(y_{-}) {}
 S operator*(const S &rhs) const {
                                                                           - d[i].begin();
  int w_{-} = w;
                                                                      pivot(i, s);
  if (w<sub>_</sub> == -1) w<sub>_</sub> = rhs.w;
  assert(w_! = -1 \text{ and } w_ == rhs.w);
  return { MOD, w_,
                                                                    if (!phase(0)) return VD(n, inf);
   (x * rhs.x + y * rhs.y % MOD * w) % MOD,
(x * rhs.y + y * rhs.x) % MOD };
                                                                    VD x(n);
                                                                    for (int i = 0; i < m; ++i)</pre>
                                                                     if (p[i] < n) x[p[i]] = d[i][n + 1];
                                                                    return x;
}:
                                                                  }}
int get_root(int n, int P) {
  if (P == 2 or n == 0) return n;
                                                                  5.22 Simplex Construction
  if (qpow(n, (P - 1) / 2, P) != 1) return -1;
  auto check = [&](int x) {
                                                                  Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that for all 1~\leq~j~\leq~m,
                                                                  \sum_{1 \le i \le n} A_{ji} x_i \le b_j and x_i \ge 0 for all 1 \le i \le n.
    return qpow(x, (P - 1) / 2, P); };
  if (check(n) == P-1) return -1
  int64_t a; int w; mt19937 rnd(7122);
                                                                     1. In case of minimization, let c_i' = -c_i
  do { a = rnd() % P;
w = ((a * a - n) % P + P) % P;
                                                                     2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
  } while (check(w) != P - 1);
                                                                     3. \sum_{1 < i < n} A_{ji} x_i = b_j
  return qpow(S(P, w, a, 1), (P + 1) / 2).x;
                                                                           • \sum_{1 \le i \le n} A_{ji} x_i \le b_j
5.21 Simplex
                                                                           • \sum_{1 \le i \le n} A_{ji} x_i \ge b_j
namespace simplex {
```

4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i - x_i'$ 

## 6 Geometry

## 6.1 Basic Geometry

```
#define IM imag
#define RE real
using lld = int64_t;
using 11f = long double;
using PT = std::complex<lld>;
using PTF = std::complex<llf>;
using P = PT;
auto toPTF(PT p) { return PTF{RE(p), IM(p)}; }
int sgn(11d x) \{ return (x > 0) - (x < 0); \}
11d dot(P a, P b) { return RE(conj(a) * b); }
lld cross(P'a, P'b) { return \int M(conj(a) * b); }
int ori(P a, P b, P c) {
return sgn(cross(b - a, c - a));
bool operator<(const P &a, const P &b) {</pre>
return RE(a) != RE(b) ? RE(a) < RE(b) : IM(a) < IM(b);
int quad(P p) {
return (IM(p) == 0) // use sgn for PTF
  ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
int argCmp(P a, P b) {
 // -1 / 0 / 1 <-> < / == / > (atan2)
 int qa = quad(a), qb = quad(b);
 if (qa != qb) return sgn(qa - qb);
 return sgn(cross(b, a));
template <typename V> llf area(const V & pt) {
 11d ret = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
 ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 return ret / 2.0;
P rot90(P p) { return P{-IM(p), RE(p)}; }
PTF project(PTF p, PTF q) { // p onto q return dot(p, q) * q / dot(q, q);
11f FMOD(11f x) {
if (x < -PI) x += PI * 2;
 if (x > PI) x -= PI * 2;
return x:
```

## 6.2 Segment & Line Intersection

```
struct Segment { // closed segment
PT st, dir; // represent st + t*dir for 0<=t<=1
 Segment(PT s, PT e) : st(s), dir(e - s) {}
 static bool valid(lld p, lld q) 
  // is there t s.t. 0 <= t <= 1 && qt == p ?
  if (q < 0) q = -q, p = -p;
  return 0 <= p && p <= q;
 vector<PT> ends() const { return { st, st + dir }; }
template <typename T> bool isInter(T A, PT P) {
 if (A.dir == PT(0)) return P == A.st; // BE CAREFUL
 return cross(P - A.st, A.dir) == 0 &&
  T::valid(dot(P - A.st, A.dir), norm(A.dir));
template <typename U, typename V>
bool isInter(U A, V B) {
  if (cross(A.dir, B.dir) == 0) { // BE CAREFUL
 bool res = false;
  for (PT P: A.ends()) res |= isInter(B, P);
  for (PT P: B.ends()) res |= isInter(A, P);
 return res:
 PT D = B.st - A.st;
 11d C = cross(A.dir, B.dir);
 return U::valid(cross(D, B.dir), C) &&
  V::valid(cross(D, A.dir), C);
struct Line {
PT st, ed, dir;
Line (PT s, PT e)
  : st(s), ed(e), dir(e - s) {}
PTF intersect(const Line &A, const Line &B) {
11f t = cross(B.st - A.st, B.dir) /
```

```
llf(cross(A.dir, B.dir));
return toPTF(A.st) + PTF(t) * toPTF(A.dir);
}
```

### 6.3 2D Convex Hull

```
void make_hull(vector<pll> &dots) { // n=1 => ans = {}
sort(dots.begin(), dots.end());
vector<pll> ans(1, dots[0]);
for (int ct = 0; ct < 2; ++ct, reverse(ALL(dots)))
  for (int i = 1, t = SZ(ans); i < SZ(dots); i++) {
    while (SZ(ans) > t && ori(
        ans[SZ(ans) - 2], ans.back(), dots[i]) <= 0)
        ans.pop_back();
    ans.pb(dots[i]);
    }
ans.pop_back(), ans.swap(dots);
}</pre>
```

## 6.4 3D Convex Hull

```
// return the faces with pt indexes
int flag[MXN][MXN];
struct Point{
 ld x,y,z;
 Point operator * (const ld &b) const {
  return (Point) {x*b, y*b, z*b};}
 Point operator * (const Point &b) const {
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
Point ver(Point a, Point b, Point c) {
  return (b - a) * (c - a);}
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0
 REP(i,n) REP(j,n) flag[i][j] = 0;
 vector<Face> now;
 now.emplace_back(0,1,2);
 now.emplace_back(2,1,0);
 for (int i=3; i<n; i++){
  ftop++; vector<Face> next;
  REP(j, SZ(now)) {
   Face& f=now[j]; int ff = 0;
   ld d=(pt[i]-pt[f.a]).dot(
     ver(pt[f.a], pt[f.b], pt[f.c]));
   if (d <= 0) next.push_back(f);</pre>
   if (d > 0) ff=ftop;
   else if (d < 0) ff=-ftop;</pre>
   flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
  REP(j, SZ(now)) {
Face& f=now[j];
   if (flag[f.a][f.b] > 0 &&
     flag[f.a][f.b] != flag[f.b][f.a])
    next.emplace_back(f.a,f.b,i);
   if (flag[f.b][f.c] > 0 &&
     flag[f.b][f.c] != flag[f.c][f.b])
    next.emplace_back(f.b,f.c,i);
   if (flag[f.c][f.a] > 0 &&
     flag[f.c][f.a] != flag[f.a][f.c])
    next.emplace_back(f.c,f.a,i);
  now=next;
 return now;
```

## 6.5 2D Farthest Pair

```
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {
  while(abs(cross(stk[i+1]-stk[i],
    stk[(pos+1)%n]-stk[i])) >
    abs(cross(stk[i+1]-stk[i],
    stk[pos]-stk[i]))) pos = (pos+1)%n;
ans = max({ans, dis(stk[i], stk[pos]),
    dis(stk[i+1], stk[pos]));
}
```

## 6.6 kD Closest Pair (3D ver.)

```
11f solve(vector<P> v) {
 shuffle(v.begin(), v.end(), mt19937());
 unordered_map<lld, unordered_map<lld,
  unordered_map<lld, int>>> m;
 llf d = dis(v[0], v[1]);
 auto Idx = [&d] (11f x) -> 11d {
 return round(x * 2 / d) + 0.1; };
auto rebuild_m = [&m, &v, &Idx](int k) {
  m.clear();
  for (int i = 0; i < k; ++i)
   m[Idx(v[i].x)][Idx(v[i].y)]
[Idx(v[i].z)] = i;
 }; rebuild_m(2);
 for (size_t i = 2; i < v.size(); ++i) {
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),</pre>
     kz = Idx(v[i].z); bool found = false;
  for (int dx = -2; dx <= 2; ++dx) {
   const 11d nx = dx + kx;
   if (m.find(nx) == m.end()) continue;
   auto& mm = m[nx];
   for (int dy = -2; dy <= 2; ++dy) {
    const 11d ny = dy + ky;
    if (mm.find(ny) == mm.end()) continue;
    auto& mmm = mm[ny];
    for (int dz = -2; dz <= 2; ++dz) {
     const lld nz = dz + kz;
     if (mmm.find(nz) == mmm.end()) continue;
     const int p = mmm[nz];
     if (dis(v[p], v[i]) < d) {</pre>
      d = dis(v[p], v[i]);
      found = true;
     }
  if (found) rebuild_m(i + 1);
  else m[kx][ky][kz] = i;
 return d;
```

## 6.7 Simulated Annealing

```
11f anneal() {
 mt19937 rnd_engine( seed );
 uniform_real_distribution< llf > rnd( 0, 1 );
 const llf dT = 0.001;
 // Argument p
11f S_cur = calc( p ), S_best = S_cur;
for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
  // Modify p to p_prime
const llf S_prime = calc( p_prime );
  const llf delta_c = S_prime - S_cur
  llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
  if ( rnd( rnd_engine ) <= prob )</pre>
   S_cur = S_prime, p = p_prime;
  if ( S_prime < S_best ) // find min</pre>
   S_best = S_prime, p_best = p_prime;
return S_best;
```

## 6.8 Half Plane Intersection

```
/ cross(pt-line.st, line.dir)<=0 <-> pt in half plane
bool operator<(const Line &lhs, const Line &rhs) {</pre>
 if (int cmp = argCmp(lhs.dir, rhs.dir))
    return cmp == -1;
  return ori(lhs.st, lhs.ed, rhs.st) < 0;
// intersect function is in "Segment Intersect"
11f HPI(vector<Line> &lines) {
 sort(lines.begin(), lines.end());
 deque<Line> que;
 deque<PTF> pt;
  que.push_back(lines[0]);
  for (int i = 1; i < (int)lines.size(); i++) {</pre>
   if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
     continue;
#define POP(L, R) \
   while (pt.size() > 0 \
```

```
&& ori(L.st, L.ed, pt.back()) < 0) \
      pt.pop_back(), que.pop_back(); \
    while (pt.size() > 0 \
      && ori(R.st, R.ed, pt.front()) < 0) \
      pt.pop_front(), que.pop_front();
    POP(lines[i], lines[i]);
    pt.push_back(intersect(que.back(), lines[i]));
    que.push_back(lines[i]);
  POP(que.front(), que.back())
  if (que.size() <= 1 ||</pre>
    argCmp(que.front().dir, que.back().dir) == 0)
    return 0:
  pt.push_back(intersect(que.front(), que.back()));
  return area(pt);
6.9 Minkowski Sum
vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
 hull(A), hull(B);
 vector<pll> C(1, A[0] + B[0]), s1, s2;
 for(int i = 0; i < SZ(A); ++i)</pre>
  s1.pb(A[(i + 1) % SZ(A)] - A[i]);
 for(int i = 0; i < SZ(B); i++)
s2.pb(B[(i + 1) % SZ(B)] - B[i]);
 for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
  if (p2 >= SZ(B)
    | | (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
   C.pb(C.back() + s1[p1++]);
  else
   C.pb(C.back() + s2[p2++]);
 return hull(C), C;
6.10 Circle Class
struct Circle { PTF o; llf r; };
vector<llf> intersectAngle(Circle A, Circle B) {
 PTF dir = B.o - A.o; llf d2 = norm(dir);
 if (norm(A.r - B.r) >= d2) // norm(x) := |x|^2
 if (A.r < B.r) return {-PI, PI}; // A in B</pre>
  else return {}; // B in A
 if (norm(A.r + B.r) <= d2) return {};</pre>
 11f dis = abs(dir), theta = arg(dir);
 11f phi = acos((A.r * A.r + d2 - B.r * B.r) /
   (2 * A.r * dis));
 11f L = FMOD(theta - phi), R = FMOD(theta + phi);
 return { L, R };
vector<PTF> intersectPoint(Circle a, Circle b) {
 11f d = abs(a.o - b.o);
 if (d >= b.r+a.r || d <= abs(b.r-a.r)) return {};</pre>
 11f dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
 PTF dir = (a.o - b.o) / d;
 PTF u = dir*d1 + b.o;
PTF v = rot90(dir) * sqrt(max<llf>(0, b.r*b.r-d1*d1));
 return {u + v, u - v};
6.11 Intersection of line and Circle
vector<PTF> line_interCircle(const PTF &p1,
  const PTF &p2, const PTF &c, const double r) {
 PTF ft = p1 + project(c-p1, p2-p1), vec = p2-p1;
 llf dis = abs(c - ft);
 if (abs(dis - r) < eps) return {ft};</pre>
 if (dis > r) return {};
 vec = vec * sqrt(r * r - dis * dis) / abs(vec);
```

```
return {ft + vec, ft - vec};
```

## 6.12 Intersection of Polygon and Circle

```
// Divides into multiple triangle, and sum up
   test by HDU2892
11f _area(PTF pa, PTF pb, llf r) {
 if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
 if (abs(pb) < eps) return 0;</pre>
 11f S, h, theta;
 llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
llf cosB = dot(pb, pb - pa) / a / c, B = acos(cosB);
 11f cosC = dot(pa, pb) / a / b, C = acos(cosC);
```

```
if (a > r) {
 S = (C / 2) * r * r;
 h = a * b * sin(C) / c;
  if (h < r && B < PI / 2)
   S = (acos(h / r) * r * r - h * sqrt(r*r - h*h));
} else if (b > r) {
  theta = PI - B - asin(sin(B) / r * a);
 S = 0.5 * a*r * sin(theta) + (C - theta) / 2 * r*r;
} else
 S = 0.5 * sin(C) * a * b;
return S;
11f area_poly_circle(const vector<PTF> &poly,
  const PTF &0, const llf r) {
11f S = 0:
for (int i = 0, N = poly.size(); i < N; ++i)</pre>
 S += _area(poly[i] - 0, poly[(i + 1) % N] - 0, r) *
  ori(0, poly[i], poly[(i + 1) % N]);
return fabs(S);
6.13 Point & Hulls Tangent
#define above(P, Vi, Vj) (ori(P, Vi, Vj) > 0) // true
    if Vi is above Vj
#define below(P, Vi, Vj) (ori(P, Vi, Vj) < 0) // true</pre>
    if Vi is below Vj
// Rtangent_PointPolyC(): binary search for convex
```

```
polygon right tangent
    Input: P = a \ 2D \ point \ (exterior \ to \ the \ polygon)
//
        n = number of polygon vertices
//
        V = array of vertices for a 2D convex polygon
    with V[n] = V[0]
   Return: index "i" of rightmost tangent point V[i]
int Rtangent_PointPolyC(PT P, int n, PT *V) {
int a, b, c;
int upA, dnC;
if (below(P, V[1], V[0]) && !above(P, V[n - 1], V[0]))
 return 0:
for (a = 0, b = n;;) {
 c = (a + b) / 2;
 dnC = below(P, V[c + 1], V[c]);
 if (dnC && !above(P, V[c - 1], V[c]))
  return c:
 upA = above(P, V[a + 1], V[a]);
 if (upA) {
  if (dnC) {
   b = c;
  } else {
   if (above(P, V[a], V[c]))
    b = c;
   else
    a = c:
  } else {
  if (!dnC) {
   a = c;
   } else {
   if (below(P, V[a], V[c]))
    b = c;
   else
    a = c:
 }
// Ltangent_PointPolyC(): binary search for convex
    polygon left tangent
   Input: P = a \ 2D \ point \ (exterior \ to \ the \ polygon)
        n = number of polygon vertices
//
        V = array of vertices for a 2D convex polygon
    with V[n]=V[0]
  Return: index "i" of leftmost tangent point V[i]
int Ltangent_PointPolyC(PT P, int n, PT *V) {
int a, b, c;
int dnA, dnC;
if (above(P, V[n - 1], V[0]) && !below(P, V[1], V[0]))
 return 0;
```

```
for (a = 0, b = n;;) {
c = (a + b) / 2;
dnC = below(P, V[c + 1], V[c]);
if (above(P, V[c - 1], V[c]) && !dnC)
  return c;
dnA = below(P, V[a + 1], V[a]);
if (dnA)
 if (!dnC) {
  b = c;
  } else {
   if (below(P, V[a], V[c]))
   b = c;
   else
   a = c;
} else {
 if (dnC) {
  a = c;
  } else {
  if (above(P, V[a], V[c]))
   b = c;
   else
   a = c:
}
```

#### 6.14 Convex Hulls Tangent

```
// RLtangent_PolyPolyC(): get the RL tangent between
    two convex polygons
    Input: m = number of vertices in polygon 1
        V = array of vertices for convex polygon 1 with
//
     V[m]=V[0]
        n = number of vertices in polygon 2
//
        W = array of vertices for convex polygon 2 with
     W[n]=W[0]
//
    Output: *t1 = index of tangent point V[t1] for
    polygon 1
        *t2 = index of tangent point W[t2] for polygon
void RLtangent_PolyPolyC(int m, PT *V, int n, PT *W,
    int *t1, int *t2) {
 int ix1, ix2; // search indices for polygons 1 and 2
 // first get the initial vertex on each polygon
 ix1 = Rtangent_PointPolyC(W[0], m, V); // right
    tangent from W[0] to V
 ix2 = Ltangent_PointPolyC(V[ix1], n, W); // left
    tangent from V[ix1] to W
 // ping-pong linear search until it stabilizes
 int done = false; // flag when done
 while (done == false) {
  done = true: // assume done until..
  while (ori(W[ix2], V[ix1], V[ix1 + 1]) <= 0) {</pre>
   ++ix1; // get Rtangent from W[ix2] to V
  while (ori(V[ix1], W[ix2], W[ix2 - 1]) >= 0) {
   --ix2;
             // get Ltangent from V[ix1] to W
   done = false; // not done if had to adjust this
  }
 *t1 = ix1;
 *t2 = ix2;
 return;
```

#### **Tangent line of Two Circle** 6.15

```
vector<Line>
tanline(const Circle &c1, const Circle &c2, int sign1){
 // sign1 = 1 for outer tang, -1 for inter tang
 vector<Line> ret;
 if (norm(c1.o - c2.o) < eps) return ret;</pre>
 11f d = abs(c1.o - c2.o);
PTF v = (c2.o - c1.o) / d;

11f c = (c1.r - sign1 * c2.r) / d;
 if (c * c > 1) return ret;
 llf h = sqrt(max<llf>(0, 1 - c * c));
 for (int sign2 : {1, -1}) {
 PTF n = c * v + sign2 * h * rot90(v);
```

tree[M].R = build\_tree(M+1, R, d+1);

if (tree[M].R) {

```
PTF p1 = c1.o + n * c1.r;
                                                                    tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
                                                                    tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
  PTF p2 = c2.o + n * (c2.r * sign1);
  if (norm(p2 - p1) < eps)
                                                                    tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
   p2 = p1 + rot90(c2.o - c1.o);
                                                                    tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
  ret.push_back({p1, p2});
                                                                   return tree+M;
 return ret;
}
                                                                  int touch(Node* r, int x, int y, LL d2){
                                                                   LL dis = sqrt(d2)+1;
6.16 Minimum Covering Circle
                                                                   if (x<r->x1-dis || x>r->x2+dis ||
template<typename P>
                                                                     y<r->y1-dis || y>r->y2+dis)
Circle getCircum(const P &a, const P &b, const P &c){
                                                                    return 0:
                                                                   return 1;
 Real a1 = a.x-b.x, b1 = a.y-b.y;
 Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
                                                                  void nearest(Node* r,int x,int y,int &mID,LL &md2) {
 Real a2 = a.x-c.x, b2 = a.y-c.y;
 Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
                                                                   if (!r || !touch(r, x, y, md2)) return;
LL d2 = dis2(r->x, r->y, x, y);
 Circle cc:
                                                                   if (d2 < md2 \mid | (d2 == md2 && mID < r->id)) {
 cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
                                                                    mID = r -> id;
 cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
 cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
                                                                    md2 = d2;
 return cc;
                                                                   // search order depends on split dim
                                                                   if ((r->f == 0 && x < r->x) ||
                                                                     (r->f == 1 \&\& y < r->y)) {
template<typename P>
                                                                    nearest(r->L, x, y, mID, md2);
nearest(r->R, x, y, mID, md2);
Circle MinCircleCover(const vector<P>& pts){
 random_shuffle(pts.begin(), pts.end());
 Circle c = { pts[0], 0 };
                                                                   } else {
 for(int i=0;i<(int)pts.size();i++){</pre>
                                                                    nearest(r->R, x, y, mID, md2);
                                                                    nearest(r\rightarrow L, x, y, mID, md2);
  if (dist(pts[i], c.o) <= c.r) continue;</pre>
  c = { pts[i], 0 };
  for (int j = 0; j < i; j++) {
  if(dist(pts[j], c.o) <= c.r) continue;</pre>
                                                                  int query(int x, int y) {
                                                                   int id = 1029384756;
   c.o = (pts[i] + pts[j]) / 2;
   c.r = dist(pts[i], c.o);
                                                                   LL d2 = 102938475612345678LL;
   for (int k = 0; k < j; k++) {
                                                                   nearest(root, x, y, id, d2);
    if (dist(pts[k], c.o) <= c.r) continue;</pre>
                                                                   return id:
    c = getCircum(pts[i], pts[j], pts[k]);
                                                                } tree;
  }
                                                                 6.18 Rotating Sweep Line
 }
                                                                 void rotatingSweepLine(pair<int, int> a[], int n) {
 return c;
                                                                  vector<pair<int, int>> 1;
1.reserve(n * (n - 1) / 2);
       KDTree (Nearest Point)
6.17
                                                                  for (int i = 0; i < n; ++i)</pre>
const int MXN = 100005;
                                                                   for (int j = i + 1; j < n; ++j)
struct KDTree {
                                                                    1.emplace_back(i, j);
 struct Node {
                                                                  sort(1.begin(), 1.end(), [&a](auto &u, auto &v){
                                                                   11d udx = a[u.first].first - a[u.second].first;
11d udy = a[u.first].second - a[u.second].second;
  int x,y,x1,y1,x2,y2;
 int id,f;
Node *L, *R;
                                                                   1ld vdx = a[v.first].first - a[v.second].first;
                                                                   11d vdy = a[v.first].second - a[v.second].second;
 } tree[MXN], *root;
 int n;
                                                                   if (udx == 0 or vdx == 0) return not udx == 0;
                                                                   int s = sgn(udx * vdx);
 LL dis2(int x1, int y1, int x2, int y2) {
 LL dx = x1-x2, dy = y1-y2;
                                                                   return udy * vdx * s < vdy * udx * s;</pre>
 return dx*dx+dy*dy;
                                                                  }):
                                                                  vector<int> idx(n), p(n);
 static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
                                                                  iota(idx.begin(), idx.end(), 0);
                                                                  sort(idx.begin(), idx.end(), [&a](int i, int j){
  return a[i] < a[j]; });</pre>
 static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 void init(vector<pair<int,int>> ip) {
                                                                  for (int i = 0; i < n; ++i) p[idx[i]] = i;</pre>
 n = ip.size();
 for (int i=0; i<n; i++) {
  tree[i].id = i;</pre>
                                                                  for (auto [i, j]: 1) {
  // do here
   tree[i].x = ip[i].first;
                                                                   swap(p[i], p[j]);
                                                                   idx[p[i]] = i, idx[p[j]] = j;
   tree[i].y = ip[i].second;
  root = build_tree(0, n-1, 0);
                                                                 6.19
                                                                        Circle Cover
 Node* build_tree(int L, int R, int d) {
                                                                 const int N = 1021;
  if (L>R) return nullptr
  int M = (L+R)/2; tree[M].f = d%2;
                                                                 struct CircleCover {
  nth_element(tree+L,tree+M,tree+R+1,d%2?cmpy:cmpx);
tree[M].x1 = tree[M].x2 = tree[M].x;
                                                                  int C;
                                                                  Cir c[N]
  tree[M].y1 = tree[M].y2 = tree[M].y;
                                                                  bool g[N][N], overlap[N][N];
  tree[M].L = build_tree(L, M-1, d+1);
                                                                  // Area[i] : area covered by at least i circles
double Area[ N ];
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
                                                                  void init(int _C){ C = _C;}
   tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
                                                                  struct Teve {
                                                                   PTF p; double ang; int add;
   tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
                                                                   Teve() {}
```

Teve(PTF \_a, double \_b, int \_c):p(\_a), ang(\_b), add(

bool operator<(const Teve &a)const

```
{return ang < a.ang;}
}eve[N * 2];
// strict: x = 0, otherwise x = -1
bool disjuct(Cir &a, Cir &b, int x)
{return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
bool contain(Cir &a, Cir &b, int x)
{return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
bool contain(int i, int j) {
 /* c[j] is non-strictly in c[i]. */
 return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c
   [j].R) == 0 \&\& i < j)) \&\& contain(c[i], c[j], -1);
void solve(){
 fill_n(Area, C + 2, 0);
 for(int i = 0; i < C; ++i)</pre>
  for(int j = 0; j < C; ++j)
  overlap[i][j] = contain(i, j);</pre>
 for(int i = 0; i < C; ++i)
  for(int j = 0; j < C; ++j)
   g[i][j] = !(overlap[i][j] || overlap[j][i] ||
  disjuct(c[i], c[j], -1));
 for(int i = 0; i < C; ++i){
  int E = 0, cnt = 1;
  for(int j = 0; j < C; ++j)</pre>
   if(j != i && overlap[j][i])
    ++cnt;
  for(int j = 0; j < C; ++j)</pre>
   if(i != j && g[i][j]) {
    auto IP = intersectPoint(c[i], c[j]);
    PTF aa = IP[0], bb = IP[1];
    llf A = arg(aa-c[i].0), B = arg(bb-c[i].0);
    eve[E++] = Teve(bb,B,1), eve[E++]=Teve(aa,A,-1);
    if(B > A) ++cnt;
  if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
   sort(eve, eve + E);
   eve[E] = eve[0];
   for(int j = 0; j < E; ++j){}
    cnt += eve[j].add;
    Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
    double theta = eve[j + 1].ang - eve[j].ang;
    if (theta < 0) theta += 2. * pi;</pre>
    Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
```

# 7 Stringology

## 7.1 Hash

```
class Hash {
  private:
    static constexpr int P = 127, Q = 1051762951;
  vector<int> h, p;
  public:
    void init(const string &s){
      h.assign(s.size()+1, 0); p.resize(s.size()+1);
      for (size_t i = 0; i < s.size(); ++i)
         h[i + 1] = add(mul(h[i], P), s[i]);
      generate(p.begin(), p.end(), [x=1, y=1, this]()
         mutable{y=x; x=mul(x, P); return y;});
    }
    int query(int l, int r){ // 1-base (l, r]
      return sub(h[r], mul(h[1], p[r-1]));}
};</pre>
```

## 7.2 Suffix Array

```
namespace sfx {
bool _t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], _c[maxn * 2];
int x[maxn], _p[maxn], _q[maxn * 2];
// sa[i]: sa[i]-th suffix is the
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *a, int *c, int n, int z) {
  memset(a, 0, sizeof(int) * n);
```

```
memcpy(x, c, sizeof(int) * z);
void induce(int *a,int *c,int *s,bool *t,int n,int z){
 memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
  if (a[i] && !t[a[i] - 1])
   a[x[s[a[i] - 1]]++] = a[i] - 1;
 memcpy(x, c, sizeof(int) * z);
 for (int i = n - 1; i >= 0; --i)
if (a[i] && t[a[i] - 1])
    a[--x[s[a[i] - 1]]] = a[i] - 1;
void sais(int *s, int *a, int *p, int *q,
bool *t, int *c, int n, int z) {
 bool uniq = t[n - 1] = true;
 int nn=0, nmxz=-1, *nsa = a+n, *ns=s+n, last=-1;
 memset(c, 0, sizeof(int) * z);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
 for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
 if (uniq)
  for (int i = 0; i < n; ++i) a[--c[s[i]]] = i;
  return;
 for (int i = n - 2; i \ge 0; --i)
  t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 pre(a, c, n, z);
for (int i = 1; i <= n - 1; ++i)</pre>
  if (t[i] && !t[i - 1])
    a[--x[s[i]]] = p[q[i] = nn++] = i;
 induce(a, c, s, t, n, z);
for (int i = 0; i < n; ++i)
  if (a[i] && t[a[i]] && !t[a[i] - 1]) {
  bool neq = last < 0 || \</pre>
   memcmp(s + a[i], s + last,
(p[q[a[i]] + 1] - a[i]) * sizeof(int));
  ns[q[last = a[i]]] = nmxz += neq;
 sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
 pre(a, c, n, z);
for (int i = nn - 1; i >= 0; --i)
  a[--x[s[p[nsa[i]]]] = p[nsa[i]];
 induce(a, c, s, t, n, z);
void build(const string &s) {
 const int n = int(s.size());
 for (int i = 0; i < n; ++i) _s[i] = s[i];
 _s[n] = 0; // s shouldn't contain 0
 sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
for(int i = 0; i < n; ++i) rev[sa[i] = sa[i+1]] = i;</pre>
 int ind = hi[0] = 0;
 for (int i = 0; i < n; ++i) {</pre>
  if (!rev[i]) {
   ind = 0;
   continue;
  while (i + ind < n && \</pre>
   s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  hi[rev[i]] = ind ? ind-- : 0;
}}
```

### 7.3 Suffix Automaton

```
struct SuffixAutomaton {
 struct node {
  int ch[K], len, fail, cnt, indeg;
  node(int L = 0) : ch{}, len(L), fail(0), cnt(0),
    indeg(0) \{ \}
 } st[N];
 int root, last, tot;
 void extend(int c) {
  int cur = ++tot;
  st[cur] = node(st[last].len + 1);
  while (last && !st[last].ch[c]) {
    st[last].ch[c] = cur;
    last = st[last].fail;
  if (!last) {
    st[cur].fail = root;
  } else {
    int q = st[last].ch[c];
    if (st[q].len == st[last].len + 1) {
      st[cur].fail = q;
```

return z;

```
7.6 Manacher
    } else {
      int clone = ++tot;
                                                                int z[maxn];
      st[clone] = st[q];
                                                                int manacher(const string& s) {
      st[clone].len = st[last].len + 1;
                                                                 string t =
      st[st[cur].fail = st[q].fail = clone].cnt = 0;
                                                                 for(char c: s) t += c, t += '.';
      while (last && st[last].ch[c] == q) {
                                                                 int 1 = 0, r = 0, ans = 0;
        st[last].ch[c] = clone;
                                                                 for (int i = 1; i < t.length(); ++i) {</pre>
        last = st[last].fail;
                                                                  z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
                                                                  while (i - z[i] \ge 0 \&\& i + z[i] < t.length()) {
    }
                                                                   if(t[i - z[i]] == t[i + z[i]]) ++z[i];
  }
                                                                   else break;
  st[last = cur].cnt += 1;
 }
                                                                  if (i + z[i] > r) r = i + z[i], l = i;
 void init(const char* s) {
 root = last = tot = 1;
                                                                 for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);</pre>
  st[root] = node(0);
                                                                 return ans:
  for (char c; c = *s; ++s) extend(c - 'a');
 int q[N];
                                                                7.7 Lexico Smallest Rotation
 void dp() {
                                                                string mcp(string s) {
  for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
                                                                 int n = s.length();
                                                                 s += s; int i = 0, j = 1;
  int head = 0, tail = 0;
                                                                 while (i < n && j < n) {</pre>
  for (int i = 1; i <= tot; i++)</pre>
                                                                  int k = 0;
    if (st[i].indeg == 0) q[tail++] = i;
                                                                  while (k < n \&\& s[i + k] == s[j + k]) k++;
  while (head != tail) {
                                                                  ((s[i+k] \leftarrow s[j+k]) ? j : i) += k + 1;
    int now = q[head++];
                                                                  j += (i == j);
    if (int f = st[now].fail) {
      st[f].cnt += st[now].cnt;
                                                                 return s.substr(i < n ? i : j, n);</pre>
      if (--st[f].indeg == 0) q[tail++] = f;
  }
                                                                7.8 Main Lorentz
                                                                vector<tuple<tuple<size_t, size_t, int, int>>> reps;
 int run(const char* s) {
                                                                void find_repetitions(const string &s, int shift = 0) {
 int now = root;
                                                                 if (s.size() <= 1)
  for (char c; c = *s; ++s)
    if (!st[now].ch[c -= 'a']) return 0;
                                                                  return:
                                                                 const size_t nu = s.size() / 2, nv = s.size() - nu;
    now = st[now].ch[c];
                                                                 string u = s.substr(0, nu), v = s.substr(nu);
                                                                 string ru(u.rbegin(), u.rend());
string rv(v.rbegin(), v.rend());
  return st[now].cnt;
                                                                 find_repetitions(u, shift);
} SAM;
                                                                 find_repetitions(v, shift + nu);
                                                                 auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u),
    z3 = Zalgo(ru + '#' + rv), z4 = Zalgo(v);
for (size_t cntr = 0; cntr < s.size(); cntr++) {</pre>
7.4 KMP
vector<int> kmp(const string &s) {
 vector<int> f(s.size(), 0);
                                                                  size_t 1; int k1, k2;
 /* f[i] = length of the longest prefix
                                                                  if (cntr < nu) {
   (excluding s[0:i]) such that it coincides
                                                                   1 = nu - cntr;
   with the suffix of s[0:i] of the same length */
                                                                   k1 = 1 < z1.size() ? z1[1] : 0;
 /* i + 1 - f[i] is the length of the
                                                                   k2 = n + 1 - 1 < z2.size() ? z2[n + 1 - 1] : 0;
   smallest recurring period of s[0:i] */
                                                                  } else {
 int k = 0:
                                                                   1 = cntr - nu + 1;
 for (int i = 1; i < (int)s.size(); ++i) {
  while (k > 0 && s[i] != s[k]) k = f[k - 1];
                                                                   k1 = n + 1 - 1 < z3.size() ? z3[n + 1 - 1] : 0;
                                                                   k2 = 1 < z4.size() ? z4[1] : 0;
  if (s[i] == s[k]) ++k;
  f[i] = k;
                                                                  if (k1 + k2 >= 1)
                                                                   reps.emplace_back(cntr, 1, k1, k2);
 return f;
vector<int> search(const string &s, const string &t) {
 // return 0-indexed occurrence of t in s
                                                                7.9 BWT
 vector<int> f = kmp(t), r;
                                                                struct BurrowsWheeler{
 for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
                                                                #define SIGMA 26
  while(k > 0 && (k==(int)t.size() \mid \mid s[i]!=t[k]))
                                                                #define BASE 'a'
   k = f[k - 1]
                                                                 vector<int> v[ SIGMA ];
  if (s[i] == t[k]) ++k;
                                                                 void BWT(char* ori, char* res){
  if (k == (int)t.size()) r.push_back(i-t.size()+1);
                                                                  // make ori -> ori + ori
                                                                  // then build suffix array
 return res;
                                                                 void iBWT(char* ori, char* res){
                                                                  for( int i = 0 ; i < SIGMA ; i ++ )</pre>
7.5 Z value
                                                                   v[ i ].clear();
                                                                  int len = strlen( ori_);
vector<int> Zalgo(const string &s) {
                                                                  for( int i = 0 ; i < len ; i ++ )
v[ ori[i] - BASE ].push_back( i );</pre>
 vector<int> z(s.size(), s.size());
 for (int i = 1, 1 = 0, r = 0; i < z[0]; ++i) { int j = clamp(r - i, 0, z[i - 1]);
                                                                  vector<int> a:
  for (; i + j < z[0] \text{ and } s[i + j] == s[j]; ++j);
                                                                  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
                                                                   for( auto j : v[ i ] ){
  if (i + (z[i] = j) > r) r = i + z[1 = i];
                                                                    a.push_back( j );
ori[ ptr ++ ] = BASE + i;
```

```
for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
    ptr = a[ ptr ];
  res[ len ] = 0;
} bwt;
```

#### **Palindromic Tree** 7.10

```
struct palindromic_tree{
struct node{
 int next[26],f,len;
 int cnt,num,st,ed; // num = depth of fail link
 node(int 1=0):f(0),len(1),cnt(0),num(0) {
  memset(next, 0, sizeof(next)); }
vector<node> st;
vector<char> s:
int last,n;
void init(){
 st.clear();s.clear();last=1; n=0;
 st.push_back(0);st.push_back(-1);
 st[0].f=1;s.push_back(-1); }
int getFail(int x){
 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
 return x;}
void add(int c){
 s.push_back(c-='a'); ++n;
 int cur=getFail(last);
  if(!st[cur].next[c]){
  int now=st.size();
  st.push_back(st[cur].len+2);
  st[now].f=st[getFail(st[cur].f)].next[c];
  st[cur].next[c]=now;
  st[now].num=st[st[now].f].num+1;
 last=st[cur].next[c];
 ++st[last].cnt;}
void dpcnt() { // cnt = #occurence in whole str
 for (int i=st.size()-1; i >= 0; i--)
  st[st[i].f].cnt += st[i].cnt;
int size(){ return st.size()-2;}
} pt;
int main() {
string s; cin >> s; pt.init();
for (int i=0; i<SZ(s); i++) {
  int prvsz = pt.size(); pt.add(s[i]);</pre>
 if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [1,r]: s.substr(1, r-1+1)
return 0;
```

#### Misc 8

#### 8.1 Theorems

## 8.1.1 Sherman-Morrison formula

$$(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1+v^{\mathsf{T}}A^{-1}u}$$

#### 8.1.2 Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} =$ d(i),  $L_{ij} = -c$  where c is the number of edge (i, j) in  $\tilde{G}$ .

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .

#### 8.1.3 Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in \mathit{E}$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

## 8.1.4 Cayley's Formula

- Given a degree sequence  $d_1, d_2, \dots, d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .

#### 8.1.5 Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1 + d_2 + \ldots + d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all  $1 \leq k \leq n$ .

### 8.1.6 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other great-

### 8.1.7 Euler's planar graph formula

 $V - E + F = C + 1, E \le 3V - 6$ (?)

#### 8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have A #{lattice points in the interior} +  $\frac{\#\{\text{lattice points on the boundary}\}}{2} - 1$ 

### 8.1.9 Matroid Intersection

Given matroids  $M_1=(G,I_1), M_2=(G,I_2)$ , find maximum  $S\in I_1\cap I_2.$  For each iteration, build the directed graph and find a shortest path from  $\boldsymbol{s}$  to  $\boldsymbol{t}$ .

```
• s \rightarrow x : S \sqcup \{x\} \in I_1
• x \to t : S \sqcup \{x\} \in I_2
• y \to x : S \setminus \{y\} \sqcup \{x\} \in I_1 (y is in the unique circuit of S \sqcup \{x\})
• x \to y : S \setminus \{y\} \sqcup \{x\} \in I_2 (y is in the unique circuit of S \sqcup \{x\})
```

Alternate the path, and |S| will increase by 1. Let  ${\cal R}$  $\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|. \ \ \text{In each iteration,} \ |E|=\mathsf{For weighted case,} \ \mathsf{assign weight} -w(x) \ \mathsf{and} \ w(x) \ \mathsf{to} \ x \in S \ \mathsf{and} \ \mathsf{and$  $\dot{S}$  and  $x \notin \dot{S}$ , resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

#### 8.2 Bitset LCS

```
scanf("%d%d", &n, &m), u = n / 64 + 1;
for (int i = 1, c; i <= n; i++)
 scanf("%d", &c), p[c].set(i);
for (int i = 1, c; i <= m; i++) {
  scanf("%d", &c), (g = f) |= p[c];</pre>
 f.shiftLeftByOne(), f.set(0);
 ((f = g - f)^{-} = g) &= g;
printf("%d\n", f.count());
```

### 8.3 Prefix Substring LCS

```
void all_lcs(string s, string t) { // 0-base
 vector<int> h(SZ(t));
 iota(ALL(h), 0);
 for (int a = 0; a < SZ(s); ++a) {
  int v = -1;
  for (int c = 0; c < SZ(t); ++c)
if (s[a] == t[c] || h[c] < v)
    swap(h[c], v);
 // LCS(s[0, a], t[b, c]) = 
// c - b + 1 - sum([h[i] >= b] | i <= c)
  // h[i] might become -1 !!
```

## 8.4 Convex 1D/1D DP

```
struct segment {
int i, 1, r
 segment() {}
 segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
 dp[0] = 0;
 deque<segment> dq; dq.push_back(segment(0, 1, n));
 for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i);
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().1 = i + 1
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
   f(i, dq.back().1) < f(dq.back().i, dq.back().1))
    dq.pop_back();
```

```
if (dq.size()) {
 int d = 1 << 20, c = dq.back().1;</pre>
 while (d >>= 1) if (c + d <= dq.back().r)</pre>
  if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
 dq.back().r = c; seg.l = c + 1;
if (seg.1 <= n) dq.push_back(seg);</pre>
```

## ConvexHull Optimization

```
struct L {
mutable int64_t a, b, p;
bool operator<(const L &r) const { return a < r.a; }</pre>
bool operator<(int64_t x) const { return p < x; }</pre>
struct DynamicHull : multiset<L, less<>> {
static const int64_t kInf = 1e18;
bool Isect(iterator x, iterator y) {
 auto Div = [](int64_t a, int64_t b) {
   return a / b - ((a ^ b) < 0 && a % b); }
 if (y == end()) { x->p = kInf; return false; }
 if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
 else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
void Insert(int64_t a, int64_t b) {
 auto z = insert({a, b, 0}), y = z++, x = y;
 while (Isect(y, z)) z = erase(z);
 if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
 while ((y = x) != begin() && (--x)->p >= y->p)
  Isect(x, erase(y));
int64_t Query(int64_t x) {
 auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
};
```

## Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
int s = 0:
for (int i = 2; i <= n; i++)
 s = (s + m) \% i;
return s:
// died at kth
int kth(int n, int m, int k){
if (m == 1) return n-1;
for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
return k;
```

## 8.7 Tree Knapsack

```
int dp[N][K]; PII obj[N];
vector<int> G[N];
void dfs(int u, int mx){
 for(int s: G[u]) {
  if(mx < obj[s].first) continue;</pre>
  for(int i=0;i<=mx-obj[s].FF;i++)</pre>
   dp[s][i] = dp[u][i];
  dfs(s, mx - obj[s].first);
  for(int i=obj[s].FF;i<=mx;i++)</pre>
   dp[u][i] = max(dp[u][i],
    dp[s][i - obj[s].FF] + obj[s].SS);
```

## 8.8 N Queens Problem

```
vector< int > solve( int n ) {
// no solution when n=2, 3
vector< int > ret;
if ( n % 6 == 2 ) {
 for ( int i = 2 ; i <= n ; i += 2 )
  ret.push_back( i );
 ret.push_back( 3 ); ret.push_back( 1 );
 for ( int i = 7 ; i <= n ; i += 2 )
  ret.push_back( i );
 ret.push_back( 5 );
} else if ( n % 6 == 3 ) {
```

```
for ( int i = 4 ; i <= n ; i += 2 )
 ret.push_back( i );
 ret.push_back( 2 );
 for ( int i = 5 ; i <= n ; i += 2 )
ret.push_back( i );</pre>
 ret.push_back( 1 ); ret.push_back( 3 );
} else {
 for ( int i = 2 ; i <= n ; i += 2 )
 ret.push_back( i );
for ( int i = 1 ; i <= n ; i += 2 )
  ret.push_back( i );
return ret:
```

## 8.9 Stable Marriage

```
1: Initialize m \in M and w \in W to free
2: while \exists free man m who has a woman w to propose to do
         w \leftarrow \text{first woman on } m \text{'s list to whom } m \text{ has not yet proposed}
        if \exists \mbox{ some pair } (m',w) \mbox{ then }
            if w prefers m to m' then
                \dot{m}' \leftarrow \textit{free}
                 (m,w) \leftarrow \mathsf{engaged}
8:
            end if
9:
        else
10:
             (m,w) \leftarrow \mathsf{engaged}
        end if
12: end while
```

## 8.10 Binary Search On Fraction

```
struct Q {
 11 p, q;
 Q go(Q b, 11 d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and \theta <= p,q <= N
Q frac_bs(ll N) {
 Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
 assert(pred(hi));
 bool dir = 1, L = 1, H = 1;
 for (; L || H; dir = !dir) {
  11 len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
   if (Q mid = hi.go(lo, len + step);
     mid.p > N || mid.q > N || dir ^ pred(mid))
    t++;
   else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
 return dir ? hi : lo;
```