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# 1 Basic

# 1.2 Debug Macro [851d50]

```
#define all(x) begin(x), end(x)
#ifdef CKISEKI
#include <experimental/iterator>
#define safe cerr<<__PRETTY_FUNCTION__<<" line "<</pre>
      _LINE__<<" safe\n"
#define debug(a...) debug_(#a, a)
#define orange(a...) orange_(#a, a)
void debug_(const char *s, auto ...a) {
  cerr << "\e[1;32m(" << s << ") = (";</pre>
 int f = 0;
 (..., (cerr << (f++ ? ", " : "") << a));
 cerr << ")\e[0m\n";</pre>
void orange_(const char *s, auto L, auto R) {
  cerr << "\e[1;33m[ " << s << " ] = [ ";</pre>
 using namespace experimental;
 copy(L, R, make_ostream_joiner(cerr, ", "));
cerr << " ]\e[0m\n";</pre>
#else
#define safe ((void)0)
#define debug(...) safe
#define orange(...) safe
#endif
```

#### 1.3 Increase Stack [b6856c]

```
const int size = 256 << 20;
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));</pre>
```

### 1.4 Pragma Optimization [6006f6]

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,arch=skylake")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
```

#### 1.5 IO Optimization [c9494b]

```
static inline int gc() {
  constexpr int B = 1<<20; static char buf[B], *p, *q;
  if (p == q) q = (p = buf) + fread(buf, 1, B, stdin);
  return q == buf ? EOF : *p++;
}</pre>
```

# 1.6 SVG Writer [57436c]

```
class SVG {
  void p(string_view s) { o << s; }
  void p(string_view s, auto v, auto... vs) {
    auto i = s.find('$');
    o << s.substr(0, i) << v, p(s.substr(i + 1), vs...);
  }
  ofstream o; string c = "red";
  public:
    SVG(auto f, auto x1, auto y1, auto x2, auto y2) : o(f) {
    p("<svg xmlns='http://www.w3.org/2000/svg' "
        "viewBox='$ $ $'>\n"
        "<style>*{stroke-width:0.5%;}</style>\n",
        x1, -y2, x2 - x1, y2 - y1); }
```

```
~SVG() { p("</svg>\n"); }
SVG &color(string nc) { return c = nc, *this; }
void line(auto x1, auto y1, auto x2, auto y2) {
p("<line x1='$' y1='$' x2='$' y2='$' stroke='$'/>\n",
 x1, -y1, x2, -y2, c); }
void circle(auto x, auto y, auto r) {
p("<circle cx='$' cy='$' r='$' stroke='$' "</pre>
  "fill='none'/>\n", x, -y, r, c); }
void text(auto x, auto y, string s, int w = 12) {
 p("<text x='$' y='$' font-size='$px'>$</text>\n",
  x, -y, w, s);
```

#### 2 **Data Structure**

# 2.1 Dark Magic [095f25]

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using namespace __gnu_pbds;
// heap tags: paring/binary/binomial/rc_binomial/thin
template<typename T>
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>, \
                  pairing_heap_tag>;
// pbds_heap::point_iterator
  ' x = pq.push(10); pq.modify(x, 87); a.join(b);
// tree tags: rb_tree_tag/ov_tree_tag/splay_tree_tag
template<typename T>
using ordered_set = tree<T, null_type, less<T>,
   rb_tree_tag, tree_order_statistics_node_update>;
// find_by_order, order_of_key
// hash tables: cc_hash_table/gp_hash_table
```

# 2.2 Link-Cut Tree [7ce2b4]

down(stk.back());

```
template <typename Val, typename SVal> class LCT {
struct node {
 int pa, ch[2];
 bool rev;
 Val v, prod, rprod;
 SVal sv, sub, vir;
 node() : pa{0}, ch{0, 0}, rev{false}, v{}, prod{},
    rprod{}, sv{}, sub{}, vir{} {};
#define cur o[u]
#define lc cur.ch[0]
#define rc cur.ch[1]
vector<node> o;
bool is_root(int u) const {
 return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u;
bool is_rch(int u) const {
 return o[cur.pa].ch[1] == u && !is_root(u);
void down(int u) {
 if (not cur.rev) return;
 if (lc) set_rev(lc);
 if (rc) set_rev(rc);
 cur.rev = false;
void up(int u) {
 cur.prod = o[lc].prod * cur.v * o[rc].prod;
 cur.rprod = o[rc].rprod * cur.v * o[lc].rprod;
 cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
void set_rev(int u) {
 swap(lc, rc);
 swap(cur.prod, cur.rprod);
 cur.rev ^= 1;
void rotate(int u) {
 int f=cur.pa, g=o[f].pa, l=is_rch(u);
if (cur.ch[l ^ 1]) o[cur.ch[l ^ 1]].pa = f;
 if (not is_root(f)) o[g].ch[is_rch(f)] = u;
 o[f].ch[l] = cur.ch[l ^ 1];
 cur.ch[l ^ 1] = f;
 cur.pa = g, o[f].pa = u;
 up(f);
void splay(int u) {
 vector<int> stk = {u};
 while (not is_root(stk.back()))
  stk.push_back(o[stk.back()].pa);
  while (not stk.empty()) {
```

```
stk.pop_back();
  for (int f = cur.pa; not is_root(u); f = cur.pa) {
   if(!is_root(f))rotate(is_rch(u)==is_rch(f)?f:u);
  rotate(u);
 }
 up(u);
 void access(int x) {
 for (int u = x, last = 0; u; u = cur.pa) {
   splay(u);
   cur.vir = cur.vir + o[rc].sub - o[last].sub;
  rc = last; up(last = u);
  splay(x);
 int find_root(int u) {
  int la = 0;
  for (access(u); u; u = lc) down(la = u);
  return la:
 void split(int x, int y) {change_root(x);access(y);}
 void change_root(int u) { access(u); set_rev(u); }
public:
 LCT(int n = 0) : o(n + 1) {}
 int add(const Val &v = {}) {
 o.push_back(v);
  return int(o.size()) - 2;
 int add(Val &&v) {
 o.emplace_back(move(v));
  return int(o.size()) - 2;
 void set_val(int u, const Val &v) {
  splay(++u); cur.v = v; up(u);
 void set_sval(int u, const SVal &v) {
  splay(++u); cur.sv = v; up(u);
 Val query(int x, int y) {
  split(++x, ++y); return o[y].prod;
 SVal subtree(int p, int u) {
  change_root(++p); access(++u);
  return cur.vir + cur.sv;
 bool connected(int u, int v) {
 return find_root(++u) == find_root(++v); }
 void link(int x, int y) {
 change_root(++x); access(++y);
 o[y].vir = o[y].vir + o[x].sub;
  up(o[x].pa = y);
 void cut(int x, int y) {
 split(++x, ++y);
 o[y].ch[0] = o[x].pa = 0; up(y);
#undef cur
#undef lc
#undef rc
2.3 LiChao Segment Tree [b9c827]
```

```
struct L {
 int m, k, id;
 L(): id(-1) {}
 L(int a, int b, int c) : m(a), k(b), id(c) {}
 int at(int x) { return m * x + k; }
class LiChao {
private:
 int n; vector<L> nodes;
 static int lc(int x) { return 2 * x + 1; }
 static int rc(int x) { return 2 * x + 2; }
 void insert(int l, int r, int id, L ln) {
  int m = (l + r) >> 1;
  if (nodes[id].id == -1)
   return nodes[id] = ln, void();
  bool atLeft = nodes[id].at(l) < ln.at(l);</pre>
  if (nodes[id].at(m) < ln.at(m))</pre>
   atLeft ^= 1, swap(nodes[id], ln);
  if (r - l == 1) return;
```

```
if (atLeft) insert(l, m, lc(id), ln);
else insert(m, r, rc(id), ln);
}
int query(int l, int r, int id, int x) {
  int m = (l + r) >> 1, ret = 0;
  if (nodes[id].id != -1) ret = nodes[id].at(x);
  if (r - l == 1) return ret;
  if (x < m) return max(ret, query(l, m, lc(id), x));
  return max(ret, query(m, r, rc(id), x));
}
public:
LiChao(int n_) : n(n_), nodes(n * 4) {}
void insert(L ln) { insert(0, n, 0, ln); }
int query(int x) { return query(0, n, 0, x); }
};</pre>
```

# 2.4 Treap [ae576c]

```
__gnu_cxx::sfmt19937 rnd(7122); // <ext/random>
namespace Treap {
struct node {
int size, pri; node *lc, *rc, *pa;
node() : size(1), pri(rnd()), lc(0), rc(0), pa(0) {}
 void pull() {
  size = 1; pa = 0;
  if (lc) { size += lc->size; lc->pa = this; }
  if (rc) { size += rc->size; rc->pa = this; }
int SZ(node *x) { return x ? x->size : 0; }
node *merge(node *L, node *R) {
 if (not L or not R) return L ? L : R;
 if (L->pri > R->pri)
 return L->rc = merge(L->rc, R), L->pull(), L;
 else
  return R->lc = merge(L, R->lc), R->pull(), R;
void splitBySize(node *o, int k, node *&L, node *&R) {
 if (not o) L = R = 0;
 else if (int s = SZ(o->lc) + 1; s <= k)
  L=o, splitBySize(o->rc, k-s, L->rc, R), L->pull();
 else
 R=o, splitBySize(o->lc, k, L, R->lc), R->pull();
 // SZ(L) == k
int getRank(node *o) { // 1-base
 int r = SZ(o->lc) + 1;
 for (; o->pa; o = o->pa)
  if (o->pa->rc == o) r += SZ(o->pa->lc) + 1;
 return r;
} // namespace Treap
```

#### 2.5 Linear Basis [138d5d]

```
template <int BITS, typename S = int> struct Basis {
 static constexpr S MIN = numeric_limits<S>::min();
 array<pair<llu, S>, BITS> b;
 Basis() { b.fill({0, MIN}); }
 void add(llu x, S p) {
  for (int i = BITS-1; i>=0; i--) if (x >> i & 1) {
   if (b[i].first == 0) return b[i]={x, p}, void();
   if (b[i].second < p)</pre>
    swap(b[i].first, x), swap(b[i].second, p);
   x ^= b[i].first;
 }
 optional<llu> query_kth(llu v, llu k) {
  vector<pair<llu, int>> o;
  for (int i = 0; i < BITS; i++)</pre>
   if (b[i].first) o.emplace_back(b[i].first, i);
  if (k >= (1ULL << o.size())) return {};</pre>
  for (int i = int(o.size()) - 1; i >= 0; i--)
   if ((k >> i & 1) ^ (v >> o[i].second & 1))
    v ^= o[i].first;
  return v;
 Basis filter(S l) {
  Basis res = *this;
  for (int i = 0; i < BITS; i++)</pre>
   if (res.b[i].second < l) res.b[i] = {0, MIN};</pre>
  return res;
|};
```

# 2.6 Binary Search On Segtree [6c61c0]

```
// find_first = l \rightarrow minimal \times s.t. check([l, x))
  find_last = r \rightarrow maximal x s.t. check([x, r))
int find_first(int l, auto &&check) {
 if (l >= n) return n + 1;
 l += sz; push(l); Monoid sum; // identity
  while ((l & 1) == 0) l >>= 1;
  if (auto s = sum + nd[l]; check(s)) {
   while (l < sz) {
    prop(l); l = (l << 1);
    if (auto nxt = sum + nd[l]; not check(nxt))
     sum = nxt, l++;
   return l + 1 - sz;
  } else sum = s, l++;
 } while (lowbit(l) != l);
 return n + 1;
int find_last(int r, auto &&check) {
 if (r <= 0) return -1;
 r += sz; push(r - 1); Monoid sum; // identity
 do {
  while (r > 1 and (r & 1)) r >>= 1;
  if (auto s = nd[r] + sum; check(s)) {
   while (r < sz) {</pre>
    prop(r); r = (r << 1) | 1;
    if (auto nxt = nd[r] + sum; not check(nxt))
     sum = nxt, r--;
   return r - sz;
  } else sum = s;
 } while (lowbit(r) != r);
 return -1;
```

# 3 Graph

#### 3.1 2-SAT (SCC) [09167a]

```
class TwoSat { // test @ CSES Giant Pizza
 int n; vector<vector<int>> G, rG, sccs;
 vector<int> ord, idx, vis, res;
 void dfs(int u) {
  vis[u] = true;
  for (int v : G[u]) if (!vis[v]) dfs(v);
  ord.push_back(u);
 void rdfs(int u) {
  vis[u] = false; idx[u] = sccs.size() - 1;
  sccs.back().push_back(u);
  for (int v : rG[u]) if (vis[v]) rdfs(v);
public:
 TwoSat(int n_{-}): n(n_{-}), G(n), rG(n), idx(n), vis(n),
    res(n) {}
 void add_edge(int u, int v) {
  G[u].push_back(v); rG[v].push_back(u);
 void orr(int x, int y) {
  if ((x ^ y) == 1) return;
  add_edge(x ^ 1, y); add_edge(y ^ 1, x);
 bool solve() {
  for (int i = 0; i < n; ++i) if (not vis[i]) dfs(i);</pre>
  for (int u : ord | views::reverse)
   if (vis[u]) sccs.emplace_back(), rdfs(u);
  for (int i = 0; i < n; i += 2)
if (idx[i] == idx[i + 1]) return false;</pre>
  vector<bool> c(sccs.size());
  for (size_t i = 0; i < sccs.size(); ++i)</pre>
   for (int z : sccs[i])
    res[z] = c[i], c[idx[z ^ 1]] = !c[i];
  return true;
 bool get(int x) { return res[x]; }
 int get_id(int x) { return idx[x];
 int count() { return sccs.size(); }
```

# **3.2** BCC [6ac6db]

```
class BCC {
 int n, ecnt, bcnt;
 vector<vector<pair<int, int>>> g;
 vector<int> dfn, low, bcc, stk;
vector<bool> ap, bridge;
 void dfs(int u, int f) {
  dfn[u] = low[u] = dfn[f] + 1;
  int ch = 0:
  for (auto [v, t] : g[u]) if (bcc[t] == -1) {
   bcc[t] = 0; stk.push_back(t);
if (dfn[v]) {
    low[u] = min(low[u], dfn[v]);
    continue;
   ++ch, dfs(v, u);
   low[u] = min(low[u], low[v]);
if (low[v] > dfn[u]) bridge[t] = true;
   if (low[v] < dfn[u]) continue;</pre>
   ap[u] = true;
   while (not stk.empty()) {
    int o = stk.back(); stk.pop_back();
    bcc[o] = bcnt;
    if (o == t) break;
   bcnt += 1;
  ap[u] = ap[u] and (ch != 1 or u != f);
public:
 BCC(int n_{-}) : n(n_{-}), ecnt(0), bcnt(0), g(n), dfn(n),
    low(n), stk(), ap(n) {}
 void add_edge(int u, int v) {
  g[u].emplace_back(v, ecnt);
  g[v].emplace_back(u, ecnt++);
 void solve() {
  bridge.assign(ecnt, false); bcc.assign(ecnt, -1);
  for (int i = 0; i < n; ++i) if (!dfn[i]) dfs(i, i);</pre>
 int bcc_id(int x) const { return bcc[x]; }
 bool is_ap(int x) const { return ap[x]; }
 bool is_bridge(int x) const { return bridge[x]; }
};
```

#### 3.3 Round Square Tree [528440]

```
struct RST {
int n; vector<vector<int>> T;
RST(auto &G) : n(G.size()), T(n) {
  vector<int> stk, vis(n), low(n);
  auto dfs = [&](auto self, int u, int d) -> void {
  low[u] = vis[u] = d; stk.push_back(u);
   for (int v : G[u]) if (!vis[v]) {
    self(self, v, d + 1);
    if (low[v] == vis[u]) {
     int cnt = T.size(); T.emplace_back();
for (int x = -1; x != v; stk.pop_back())
      T[cnt].push_back(x = stk.back());
     T[u].push_back(cnt); // T is rooted
    } else low[u] = min(low[u], low[v]);
  } else low[u] = min(low[u], vis[v]);
  };
  for (int u = 0; u < N; u++)</pre>
  if (!vis[u]) dfs(dfs, u, 1);
} // T may be forest; after dfs, stk are the roots
}; // test @ 2020 Shanghai K
```

# 3.4 Edge TCC [5a2668]

```
vector<vector<int>> ETCC(auto &adj) {
  const int n = static_cast<int>(adj.size());
  vector<int> up(n), low(n), in, out, nx, id;
  in = out = nx = id = vector<int>(n, -1);
  int dfc = 0, cnt = 0; Dsu dsu(n);
  auto merge = [&](int u, int v) {
    dsu.join(u, v); up[u] += up[v]; };
  auto dfs = [&](auto self, int u, int p) -> void {
    in[u] = low[u] = dfc++;
    for (int v : adj[u]) if (v != u) {
        if (v == p) { p = -1; continue; }
        if (in[v] == -1) {
            self(self, v, u);
        if (nx[v] == -1 && up[v] <= 1) {
            up[u] += up[v]; low[u] = min(low[u], low[v]);
        }
}</pre>
```

```
continue;
    if (up[v] == 0) v = nx[v];
    if (low[u] > low[v])
     low[u] = low[v], swap(nx[u], v);
    for (; v != -1; v = nx[v]) merge(u, v);
   } else if (in[v] < in[u]) {</pre>
    low[u] = min(low[u], in[v]); up[u]++;
    for (int &x = nx[u]; x != -1 &&
  in[x] <= in[v] && in[v] < out[x]; x = nx[x])</pre>
     merge(u, x);
    up[u]--;
  out[u] = dfc;
 for (int i = 0; i < n; i++)</pre>
  if (in[i] == -1) dfs(dfs, i, -1);
 for (int i = 0; i < n; i++)
if (dsu.anc(i) == i) id[i] = cnt++;</pre>
 vector<vector<int>> comps(cnt);
 for (int i = 0; i < n; i++)</pre>
  comps[id[dsu.anc(i)]].push_back(i);
 return comps;
} // test @ yosupo judge
```

# **3.5** DMST [75c30d]

```
using D = int64_t;
struct E { int s, t; D w; }; // 0-base
vector<int> dmst(const vector<E> &e, int n, int root) {
 using PQ = pair<min_heap<pair<D, int>>, D>;
 auto push = [](PQ &pq, pair<D, int> v) {
  pq.first.emplace(v.first - pq.second, v.second); };
 auto top = [](const PQ &pq) -> pair<D, int> {
  auto r = pq.first.top();
  return {r.first + pq.second, r.second}; };
 auto join = [&push, &top](PQ &a, PQ &b) {
  if (a.first.size() < b.first.size()) swap(a, b);</pre>
  for (; !b.first.empty(); b.first.pop())
   push(a, top(b)); };
 vector<PQ> h(n * 2);
for (size_t i = 0; i < e.size(); ++i)</pre>
  push(h[e[i].t], {e[i].w, i});
 vector<int> a(n*2), v(n*2, -1), pa(n*2, -1), r(n*2);
 iota(a.begin(), a.end(), 0);
 auto o = [&](int x) { int y;
  for (y = x; a[y] != y; y = a[y]);
  for (int ox = x; x != y; ox = x)
   x = a[x], a[ox] = y;
  return y; };
 int pc = (v[root] = n + 1) - 1;
for (int i = 0; i < n; ++i) if (v[i] == -1)</pre>
  for (int p=i; v[p]<0||v[p]==i; p=o(e[r[p]].s)) {
  if (int q = p; v[q] == i && (p = pc++, 1)) do {
   h[q].second = -h[q].first.top().first;</pre>
    join(h[pa[q] = a[q] = p], h[q]);
   } while ((q = o(e[r[q]].s)) != p);
   for(v[p]=i;!h[p].first.empty()&&o(e[top(h[p]).second
    ].s)==p;h[p].first.pop());
   r[p] = top(h[p]).second;
 vector<int> ans;
 for (int i=pc-1;i>=0;i--) if (i!=root&&v[i]!=n) {
  for (int f = e[r[i]].t; f!=-1&&v[f]!=n; f = pa[f])
   v[f] = n;
  ans.push_back(r[i]);
 return ans; // default minimize, returns edgeid array
```

#### 3.6 Dominator Tree [ea5b7c]

```
struct Dominator {
  vector<vector<int>> g, r, rdom; int tk;
  vector<int>> dfn, rev, fa, sdom, dom, val, rp;
  Dominator(int n) : g(n), r(n), rdom(n), tk(0) {
    dfn = rev = fa = sdom = dom =
      val = rp = vector<int>(n, -1); }
  void add_edge(int x, int y) { g[x].push_back(y); }
  void dfs(int x) {
    rev[dfn[x] = tk] = x;
    fa[tk] = sdom[tk] = val[tk] = tk; tk++;
```

```
for (int u : g[x]) {
   if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
   r[dfn[u]].push_back(dfn[x]);
}
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
 if (fa[x] == x) return c ? -1 : x;
 if (int p = find(fa[x], 1); p != -1) {
   if (sdom[val[x]] > sdom[val[fa[x]]])
    val[x] = val[fa[x]];
   fa[x] = p;
   return c ? p : val[x];
  } else return c ? fa[x] : val[x];
vector<int> build(int s, int n) {
  // return the father of each node in dominator tree
 dfs(s); // p[i] = -2 \text{ if i is unreachable from s}
 for (int i = tk - 1; i >= 0; --i) {
   for (int u : r[i])
    sdom[i] = min(sdom[i], sdom[find(u)]);
   if (i) rdom[sdom[i]].push_back(i);
   for (int u : rdom[i]) {
    int p = find(u);
    dom[u] = (sdom[p] == i ? i : p);
   if (i) merge(i, rp[i]);
 }
 vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)</pre>
   if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
 for (int i = 1; i < tk; ++i)
   p[rev[i]] = rev[dom[i]];
  return p;
} // test @ yosupo judge
};
```

# 3.7 Edge Coloring [029763]

```
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
for (int i = 0; i <= N; i++)</pre>
 for (int j = 0; j <= N; j++)</pre>
   C[i][j] = G[i][j] = 0;
void solve(vector<pair<int, int>> &E, int N) {
int X[kN] = {}, a;
auto update = [&](int u) {
 for (X[u] = 1; C[u][X[u]]; X[u]++);
auto color = [&](int u, int v, int c) {
 int p = G[u][v];
 G[u][v] = G[v][u] = c;
 C[u][c] = v, C[v][c] = u;
 C[u][p] = C[v][p] = 0;
 if (p) X[u] = X[v] = p;
 else update(u), update(v);
 return p;
auto flip = [&](int u, int c1, int c2) {
 int p = C[u][c1];
 swap(C[u][c1], C[u][c2]);
 if (p) G[u][p] = G[p][u] = c2;
 if (!C[u][c1]) X[u] = c1;
 if (!C[u][c2]) X[u] = c2;
 return p;
for (int i = 1; i <= N; i++) X[i] = 1;</pre>
for (int t = 0; t < E.size(); t++) {</pre>
 auto [u, v] = E[t];
 int v0 = v, c = X[u], c0 = c, d;
 vector<pair<int, int>> L; int vst[kN] = {};
 while (!G[u][v0]) {
  L.emplace_back(v, d = X[v]);
  if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
     c = color(u, L[a].first, c);
  else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
     color(u, L[a].first, L[a].second);
   else if (vst[d]) break;
  else vst[d] = 1, v = C[u][d];
 if (!G[u][v0]) {
```

```
for (; v; v = flip(v, c, d), swap(c, d));
  if (C[u][c0]) { a = int(L.size()) - 1;
    while (--a >= 0 && L[a].second != c);
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
  } else t--;
}
}
}
```

# 3.8 Centroid Decomposition [63b2fb]

```
struct Centroid {
 using G = vector<vector<pair<int, int>>>;
 vector<vector<int64_t>> Dist;
 vector<int> Pa, Dep;
 vector<int64_t> Sub, Sub2;
 vector<int> Cnt, Cnt2;
 vector<int> vis, sz, mx, tmp;
void DfsSz(const G &g, int x) {
  vis[x] = true, sz[x] = 1, mx[x] = 0;
  for (auto [u, w] : g[x]) if (not vis[u]) {
   DfsSz(g, u); sz[x] += sz[u];
   mx[x] = max(mx[x], sz[u]);
  tmp.push_back(x);
 void DfsDist(const G &g, int x, int64_t D = 0) {
 Dist[x].push_back(D); vis[x] = true;
  for (auto [u, w] : g[x])
   if (not vis[u]) DfsDist(g, u, D + w);
 void DfsCen(const G &g, int x, int D = 0, int p = -1)
  tmp.clear(); DfsSz(g, x);
  int M = tmp.size(), C = -1;
  for (int u : tmp) {
   if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
   vis[u] = false;
  DfsDist(g, C);
  for (int u : tmp) vis[u] = false;
  Pa[C] = p, vis[C] = true, Dep[C] = D;
  for (auto [u, w] : g[C])
   if (not vis[u]) DfsCen(g, u, D + 1, C);
 Centroid(int N, G g)
   : Sub(N), Sub2(N), Cnt(N), Cnt2(N), Dist(N), Pa(N),
    Dep(N), vis(N), sz(N), mx(N) { DfsCen(g, 0); }
 void Mark(int v) {
  int x = v, z = -1;
  for (int i = Dep[v]; i >= 0; --i) {
   Sub[x] += Dist[v][i], Cnt[x]++;
   if (z != -1)
    Sub2[z] += Dist[v][i], Cnt2[z]++;
   x = Pa[z = x];
  }
 int64_t Query(int v) {
  int64_t res = 0;
  int x = v, z = -1;
  for (int i = Dep[v]; i >= 0; --i) {
   res += Sub[x] + 1LL * Cnt[x] * Dist[v][i];
   if (z != -1)
    res -= Sub2[z] + 1LL * Cnt2[z] * Dist[v][i];
  x = Pa[z = x];
  return res;
}
```

#### 3.9 Lowbit Decomposition [760ac1]

```
class LBD {
  int timer, chains;
  vector<vector<int>> G;
  vector<iint>> tl, tr, chain, head, dep, pa;
  // chains : number of chain
  // tl, tr[u] : subtree interval in the seq. of u
  // head[i] : head of the chain i
  // chian[u] : chain id of the chain u is on
  void predfs(int u, int f) {
  dep[u] = dep[pa[u] = f] + 1;
  for (int v : G[u]) if (v != f) {
    predfs(v, u);
}
```

```
National Taiwan University - ckiseki
   if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
    chain[u] = chain[v];
  if (chain[u] == 0) chain[u] = ++chains;
 void dfschain(int u, int f) {
  tl[u] = timer++;
  if (head[chain[u]] == -1)
   head[chain[u]] = u;
  for (int v : G[u])
   if (v != f and chain[v] == chain[u])
    dfschain(v, u);
  for (int v : G[u])
  if (v != f and chain[v] != chain[u])
    dfschain(v, u);
  tr[u] = timer;
public:
 LBD(int n): timer(0), chains(0), G(n), tl(n), tr(n),
    chain(n), head(n + 1, -1), dep(n), pa(n) {}
 void add_edge(int u, int v) {
  G[u].push_back(v); G[v].push_back(u);
 void decompose() { predfs(0, 0); dfschain(0, 0); }
 PII get_subtree(int u) { return {tl[u], tr[u]}; }
 vector<PII> get_path(int u, int v) {
  vector<PII> res;
  while (chain[u] != chain[v]) {
   if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
   int s = head[chain[u]];
   res.emplace_back(tl[s], tl[u] + 1);
   u = pa[s];
  if (dep[u] < dep[v]) swap(u, v);</pre>
  res.emplace_back(tl[v], tl[u] + 1);
  return res:
 }
};
3.10 Virtual Tree [ad5cf5]
vector<pair<int, int>> build(vector<int> vs, int r) {
 vector<pair<int, int>> res;
 sort(vs.begin(), vs.end(), [](int i, int j) {
  return dfn[i] < dfn[j]; });</pre>
 vector<int> s = {r};
for (int v : vs) if (v != r) {
```

```
vector<pair<int, int>> build(vector<int> vs, int r) {
  vector<pair<int, int>> res;
  sort(vs.begin(), vs.end(), [](int i, int j) {
    return dfn[i] < dfn[j]; });
  vector<int> s = {r};
  for (int v : vs) if (v != r) {
    if (int o = lca(v, s.back()); o != s.back()) {
      while (s.size() >= 2) {
      if (dfn[s[s.size() - 2]] < dfn[o]) break;
      res.emplace_back(s[s.size() - 2], s.back());
      s.pop_back();
    }
    if (s.back() != o) {
      res.emplace_back(o, s.back());
      s.back() = o;
    }
    }
    s.push_back(v);
}
for (size_t i = 1; i < s.size(); ++i)
    res.emplace_back(s[i - 1], s[i]);
    return res; // (x, y): x->y
}
```

# 3.11 Tree Hashing [707efa]

```
llu F(llu z) { // xorshift64star from iwiwi
z ^= z >> 12; z ^= z << 25; z ^= z >> 27;
return z * 2685821657736338717LL;
}
llu hsah(int u, int f) {
    llu r = 127; // bigger?
    for (int v : G[u]) if (v != f) r += F( hsah(v, u) );
    return F(r);
} // test @ UOJ 763
```

# 3.12 Mo's Algorithm on Tree

```
dfs u:
push u
iterate subtree
push u
```

```
Let P = LCA(u, v) with St(u)<=St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
3.13 Count Cycles [c7e8f2]</pre>
```

```
// ord = sort by deg decreasing, rk[ord[i]] = i

// D[i] = edge point from rk small to rk big

for (int x : ord) { // c3

for (int y : D[x]) vis[y] = 1;

for (int y : D[x]) for (int z : D[y]) c3 += vis[z];

for (int y : D[x]) vis[y] = 0;

}

for (int x : ord) { // c4

for (int x : ord) { // c4

for (int y : D[x]) for (int z : adj[y])

if (rk[z] > rk[x]) c4 += vis[z]++;

for (int y : D[x]) for (int z : adj[y])

if (rk[z] > rk[x]) --vis[z];

} // both are O(M*sqrt(M)), test @ 2022 CCPC guangzhou
```

# 3.14 MaximalClique [293730]

```
// contain a self loop u to u, than u won't in clique
template <size_t maxn> class MaxClique {
private:
 using bits = bitset<maxn>;
 bits popped, G[maxn], ans;
 size_t deg[maxn], deo[maxn], n;
 void sort_by_degree() {
  popped.reset();
  for (size_t i = 0; i < n; ++i)</pre>
   deg[i] = G[i].count();
  for (size_t i = 0; i < n; ++i) {</pre>
   size_t mi = maxn, id = 0;
   for (size_t j = 0; j < n; ++j)
  if (not popped[j] and deg[j] < mi)</pre>
     mi = deg[id = j];
   popped[deo[i] = id] = 1;
   for (size_t u = G[i]._Find_first(); u < n;</pre>
     u = G[i]._Find_next(u))
     --deg[u];
 void BK(bits R, bits P, bits X) {
  if (R.count() + P.count() <= ans.count()) return;</pre>
  if (not P.count() and not X.count()) {
   if (R.count() > ans.count()) ans = R;
   return;
  /* greedily chosse max degree as pivot
  bits cur = P \mid X; size_t pivot = 0, sz = 0;
  for ( size_t u = cur._Find_first() ;
   u < n ; u = cur.\_Find\_next(u)
    if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
  cur = P & ( ~G[ pivot ] );
  */ // or simply choose first
  bits cur = P & (~G[(P | X)._Find_first()]);
  for (size_t u = cur._Find_first(); u < n;</pre>
    u = cur._Find_next(u)) {
   if (R[u]) continue;
   R[u] = 1;
   BK(R, P & G[u], X & G[u]);
   R[u] = P[u] = 0, X[u] = 1;
public:
 void init(size_t n_) {
  n = n_{;}
  for (size_t i = 0; i < n; ++i) G[i].reset();</pre>
  ans.reset();
 void add_edges(int u, bits S) { G[u] = S; }
 void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
 int solve() {
  sort_by_degree(); // or simply iota( deo... )
  for (size_t i = 0; i < n; ++i)</pre>
   deg[i] = G[i].count();
  bits pob, nob = 0; pob.set();
  for (size_t i = n; i < maxn; ++i) pob[i] = 0;
for (size_t i = 0; i < n; ++i) {</pre>
   size_t v = deo[i];
   bits tmp;
   tmp[v] = 1;
   BK(tmp, pob & G[v], nob & G[v]);
```

```
pob[v] = 0, nob[v] = 1;
}
return static_cast<int>(ans.count());
}
};
```

### 3.15 MaximumClique [aee5d8]

```
constexpr size_t kN = 150; using bits = bitset<kN>;
struct MaxClique {
bits G[kN], cs[kN];
int ans, sol[kN], q, cur[kN], d[kN], n;
void init(int _n) {
 n = _n;
  for (int i = 0; i < n; ++i) G[i].reset();</pre>
void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
void pre_dfs(vector<int> &v, int i, bits mask) {
 if (i < 4) {
  for (int x : v) d[x] = (int)(G[x] \& mask).count();
  sort(all(v), [&](int x, int y) {
    return d[x] > d[y]; });
 vector<int> c(v.size());
  cs[1].reset(), cs[2].reset();
  int l = max(ans - q + 1, 1), r = 2, tp = 0, k;
  for (int p : v) {
   for (k = 1; (cs[k] & G[p]).any(); ++k);
   if (k >= r) cs[++r].reset();
   cs[k][p] = 1;
   if (k < l) v[tp++] = p;
  for (k = l; k < r; ++k)
   for (auto p = cs[k]._Find_first();
    p < kN; p = cs[k]._Find_next(p))
   v[tp] = (int)p, c[tp] = k, ++tp;
 dfs(v, c, i + 1, mask);
void dfs(vector<int> &v, vector<int> &c,
   int i, bits mask) {
 while (!v.empty()) {
   int p = v.back(); v.pop_back(); mask[p] = 0;
   if (q + c.back() <= ans) return;</pre>
   cur[q++] = p;
   vector<int> nr;
   for (int x : v) if (G[p][x]) nr.push_back(x);
   if (!nr.empty()) pre_dfs(nr, i, mask & G[p]);
   else if (q > ans) ans = q, copy_n(cur, q, sol);
   c.pop_back(); --q;
int solve() {
 vector<int> v(n); iota(all(v), 0);
  ans = q = 0; pre_dfs(v, 0, bits(string(n, '1')));
  return ans; // sol[0 ~ ans-1]
} cliq; // test @ yosupo judge
```

# 3.16 Minimum Mean Cycle [e23bc0]

```
// WARNING: TYPE matters
struct Edge { int s, t; llf c; };
llf solve(vector<Edge> &e, int n) {
  // O(VE), returns inf if no cycle, mmc otherwise
 vector<VI> prv(n + 1, VI(n)), prve = prv;
 vector<vector<llf>>> d(n + 1, vector<llf>(n, inf));
 d[0] = vector<llf>(n, 0);
 for (int i = 0; i < n; i++) {</pre>
  for (int j = 0; j < (int)e.size(); j++) {</pre>
   auto [s, t, c] = e[j];
   if (d[i][s] < inf && d[i + 1][t] > d[i][s] + c) {
    d[i + 1][t] = d[i][s] + c;
    prv[i + 1][t] = s; prve[i + 1][t] = j;
  }
 llf mmc = inf; int st = -1;
 for (int i = 0; i < n; i++) {</pre>
  llf avg = -inf;
  for (int k = 0; k < n; k++) {
   if (d[n][i] < inf - eps)
    avg = max(avg, (d[n][i] - d[k][i]) / (n - k));
   else avg = inf;
```

```
if (avg < mmc) tie(mmc, st) = tie(avg, i);
}
if (st == -1) return inf;
vector<int> vst(n), eid, cycle, rho;
for (int i = n; !vst[st]; st = prv[i--][st]) {
  vst[st]++; eid.emplace_back(prve[i][st]);
  rho.emplace_back(st);
}
while (vst[st] != 2) {
  int v = rho.back(); rho.pop_back();
  cycle.emplace_back(v); vst[v]++;
}
reverse(all(eid)); eid.resize(cycle.size());
return mmc;
}
```

# 4 Flow & Matching 4.1 HopcroftKarp [6fd530]

```
vector<int> l, r, a, p; int ans; queue<int> q;
 HK(int n, int m, auto \&g) : l(n,-1),r(m,-1),ans(0) {
  do {
   for (int i = 0; i < n; i++)</pre>
   if (l[i] == -1) q.push(a[i] = p[i] = i);
 } while (bfs(g));
 bool bfs(auto &g) {
  // bitset<maxn> nvis, t; nvis.set();
  for (int z, x; !q.empty(); q.pop())
   // or use _Find_first and _Find_next here
   if (l[a[x = q.front()]] == -1) for (int y: g[x]) {
    // nvis.reset(v);
    if (r[y] == -1) {
    for (z = y; z != -1;)
     r[z] = x, swap(l[x], z), x = p[x];
     return ++ans, true;
   } else if (p[r[y]] == -1)
    q.push(z = r[y]), p[z] = x, a[z] = a[x];
  return false;
 }
};
```

### 4.2 Dijkstra Cost Flow [06d5f2]

```
template <typename F, typename C> class MCMF {
 static constexpr F INF_F = numeric_limits<F>::max();
 static constexpr C INF_C = numeric_limits<C>::max();
 struct E {
  int to, r; F f; C c;
  E(int a, int b, F x, C y)
   : to(a), r(b), f(x), c(y) {}
 vector<vector<E>> g; vector<pair<int, int>> f;
 vector<F> up; vector<C> d, h;
 optional<pair<F, C>> step(int S, int T) {
  priority_queue<pair<C, int>> q;
  q.emplace(d[S] = 0, S), up[S] = INF_F;
  while (not q.empty()) {
   auto [l, u] = q.top(); q.pop();
if (up[u] == 0 or l != -d[u]) continue;
for (int i = 0; i < int(g[u].size()); ++i) {</pre>
    auto e = g[u][i]; int v = e.to;
    auto nd = d[u] + e.c + h[u] - h[v];
    if (e.f <= 0 or d[v] <= nd) continue;</pre>
    f[v] = \{u, i\}; up[v] = min(up[u], e.f);
    q.emplace(-(d[v] = nd), v);
  if (d[T] == INF_C) return nullopt;
  for (size_t i = 0; i < d.size(); i++) h[i]+=d[i];</pre>
  for (int i = T; i != S; i = f[i].first) {
   auto &eg = g[f[i].first][f[i].second];
   eg.f -= up[T]; g[eg.to][eg.r].f += up[T];
  return pair{up[T], h[T]};
 }
public:
 MCMF(int n) : g(n), f(n), up(n), d(n, INF_C) {}
 void add_edge(int s, int t, F c, C w) {
  g[s].emplace_back(t, int(g[t].size()), c, w);
  g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
```

```
pair<F, C> solve(int a, int b) {
 h.assign(g.size(), 0);
  F c = 0; C w = 0;
 while (auto r = step(a, b)) {
   c += r->first, w += r->first * r->second;
   fill(d.begin(), d.end(), INF_C);
  return {c, w};
}
};
```

# 4.3 Dinic [659ddd]

```
template <typename Cap = int64_t> class Dinic {
private:
 struct E { int to, rev; Cap cap; }; int n, st, ed;
 vector<vector<E>> G; vector<size_t> lv, idx;
 bool BFS() {
  lv.assign(n, 0); idx.assign(n, 0);
  queue<int> bfs; bfs.push(st); lv[st] = 1;
  while (not bfs.empty()) {
   int u = bfs.front(); bfs.pop();
for (auto e: G[u]) if (e.cap > 0 and !lv[e.to])
    bfs.push(e.to), lv[e.to] = lv[u] + 1;
  return lv[ed];
 Cap DFS(int u, Cap f = numeric_limits<Cap>::max()) {
  if (u == ed) return f;
  Cap ret = 0;
  for (auto &i = idx[u]; i < G[u].size(); ++i) {</pre>
   auto &[to, rev, cap] = G[u][i];
   if (cap <= 0 or lv[to] != lv[u] + 1) continue;</pre>
   Cap nf = DFS(to, min(f, cap));
   ret += nf; cap -= nf; f -= nf;
   G[to][rev].cap += nf;
   if (f == 0) return ret;
  if (ret == 0) lv[u] = 0;
  return ret;
public:
 void init(int n_) { G.assign(n = n_, vector<E>()); }
 void add_edge(int u, int v, Cap c) {
   G[u].push_back({v, int(G[v].size()), c});
   G[v].push_back({u, int(G[u].size())-1, 0});
 Cap max_flow(int st_, int ed_) {
  st = st_, ed = ed_; Cap ret = 0;
  while (BFS()) ret += DFS(st);
  return ret;
}; // test @ luogu P3376
```

#### 4.4 Flow Models

- · Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - For each edge (x,y,l,u), connect  $x \to y$  with capacity u-l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v) > 0, connect  $S \to v$  with capacity in(v), otherwise, con- $\text{nect } v \to T \text{ with capacity } -in(v).$ 
    - To maximize, connect  $t\, o\,s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$  , there's no solution. Otherwise, the maximum flow from  $\boldsymbol{s}$  to  $\boldsymbol{t}$  is the answer. Also,  $\boldsymbol{f}$ is a mincost valid flow. – To minimize, let f be the maximum flow from S to T. Con-
    - nect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- · Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
  - 1. Redirect every edge:  $y \to x$  if  $(x,y) \in M$ ,  $x \to y$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited. 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source  ${\cal S}$  and sink  ${\cal T}$
  - 2. For each edge (x,y,c), connect  $x \to y$  with (cost,cap)=(c,1) if c>0, otherwise connect  $y \to x$  with (cost,cap)=(-c,1)

- 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
- 4. For each vertex v with d(v)>0, connect  $S\to v$  with (cost,cap)=0(0, d(v))
- 5. For each vertex v with d(v) < 0, connect  $v \rightarrow T$  with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let K be the sum of all weights
  - 3. Connect source  $s \to v$ ,  $v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
  - 5. For  $v \in G$ , connect it with sink  $v \to t$  with capacity K + 2T - $(\sum_{e \in E(v)} w(e)) - 2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with
  - weight w(u,v). 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 3. Find the minimum weight perfect matching on G'.
- · Submodular functions minimization
  - For a function  $f: 2^V \to \mathbb{R}$ , f is a submodular function iff

```
* \forall S, T \subseteq V, f(S) + f(T) \ge f(S \cup T) + f(S \cap T), or
```

- $* \ \forall X \subseteq Y \subseteq V, x \not\in Y, f(X \cup \{x\}) f(X) \geq f(Y \cup \{x\}) f(Y).$
- To  $\sum_{i} \theta_{i}(x_{i})$  + minimize  $\sum_{i < j} \phi_{ij}(x_i, x_j)$  $\sum_{i < j < k} \psi_{ijk}(x_i, x_j, x_k)$
- If  $\theta_i(1) \geq \theta_i(0)$ , add edge (S, i,  $\theta_i(1) \theta_i(0)$ ) and  $\theta_i(0)$  to answer; otherwise, (i, T,  $\theta_i(0) - \theta_i(1)$ ) and  $\theta_i(1)$ .
- Add edges  $(i, j, \phi_{ij}(0, 1) + \phi_{ij}(1, 0) \phi_{ij}(0, 0) \phi_{ij}(1, 1))$ .
- Denote  $x_{ijk}$  as helper nodes. Let  $P = \psi_{ijk}(0,0,0)$  + benote  $x_{ijk}$  as helper holdes. Let  $T = \psi_{ijk}(0,0,0) + \psi_{ijk}(0,1,1) + \psi_{ijk}(1,0,1) + \psi_{ijk}(1,0,0) - \psi_{ijk}(0,0,1) - \psi_{ijk}(0,1,0) - \psi_{ijk}(1,0,0) - \psi_{ijk}(1,1,1)$ . Add -P to answer. If  $P \geq 0$ , add edges  $(i,x_{ijk},P)$ ,  $(j,x_{ijk},P)$ ,  $(k,x_{ijk},P)$ ,  $(x_{ijk},T,P)$ ; otherwise  $(x_{ijk},i,-P)$ ,  $(x_{ijk},j,-P)$ ,  $(x_{ijk},k,-P)$ ,  $(S,x_{ijk},-P)$ .
- The minimum cut of this graph will be the the minimum value of the function above.

# 4.5 General Graph Matching [5f2293]

struct Matching {

```
queue<int> q; int ans, n;
vector<int> fa, s, v, pre, match;
int Find(int u) {
 return u == fa[u] ? u : fa[u] = Find(fa[u]); }
int LCA(int x, int y) {
 static int tk = 0; tk++; x = Find(x); y = Find(y);
 for (;; swap(x, y)) if (x != n) {
  if (v[x] == tk) return x;
  v[x] = tk;
  x = Find(pre[match[x]]);
 }
}
void Blossom(int x, int y, int l) {
 for (; Find(x) != l; x = pre[y]) {
  pre[x] = y, y = match[x];
  if (s[y] == 1) q.push(y), s[y] = 0;
  for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
 }
bool Bfs(auto &&g, int r) {
  iota(all(fa), 0); ranges::fill(s, -1);
 q = queue<int>(); q.push(r); s[r] = 0;
 for (; !q.empty(); q.pop()) {
  for (int x = q.front(); int u : g[x])
   if (s[u] == -1) {
    if (pre[u] = x, s[u] = 1, match[u] == n) {
     for (int a = u, b = x, last;
  b != n; a = last, b = pre[a])
      last = match[b], match[b] = a, match[a] = b;
     return true;
    q.push(match[u]); s[match[u]] = 0;
   } else if (!s[u] && Find(u) != Find(x)) {
    int l = LCA(u, x);
Blossom(x, u, l); Blossom(u, x, l);
   }
 }
 return false;
Matching(auto &&g) : ans(0), n(int(g.size())),
fa(n+1), s(n+1), v(n+1), pre(n+1, n), match(n+1, n) {
```

if (chmin(slk[x], hl[x] + hr[y] - w[x][y]))
if (pre[x] = y, !check(x)) return;

lld d = ranges::min(slk);

```
for (int x = 0; x < n; ++x)
                                                                 for (int x = 0; x < n; ++x)
                                                                  vl[x] ? hl[x] += d : slk[x] -= d;
  if (match[x] == n) ans += Bfs(g, x);
                                                                 for (int x = 0; x < n; ++x) if (vr[x]) hr[x] -= d;
} // match[x] == n means not matched
                                                                 for (int x = 0; x < n; ++x) if (!check(x)) return;
}; // test @ yosupo judge
4.6 Global Min-Cut [1f0306]
const int maxn = 500 + 5;
                                                               KM(int n_{-}, const auto \&w) : n(n_{-}), ans(0),
                                                               hl(n), hr(n), fl(n, -1), fr(fl), pre(n), q(n) {
int w[maxn][maxn], g[maxn];
                                                                for (int i = 0; i < n; ++i) hl[i]=ranges::max(w[i]);
for (int i = 0; i < n; ++i) bfs(w, i);
for (int i = 0; i < n; ++i) ans += w[i][fl[i]];</pre>
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
                                                               }
pair<int, int> phase(int n) {
                                                             };
memset(v, false, sizeof(v));
                                                                     Minimum Cost Circulation [0f0e85]
memset(g, 0, sizeof(g));
int s = -1, t = -1;
                                                              int vis[N], visc, fa[N], fae[N], head[N], mlc = 1;
while (true) {
                                                              struct ep {
 int c = -1;
  for (int i = 0; i < n; ++i) {</pre>
                                                               int to, next;
                                                               ll flow, cost;
   if (del[i] || v[i]) continue;
   if (c == -1 || g[i] > g[c]) c = i;
                                                              e[M << 1];
                                                              void adde(int u, int v, ll fl, int cs) {
                                                               e[++mlc] = {v, head[u], fl, cs};
  if (c == -1) break;
                                                               head[u] = mlc;
 v[s = t, t = c] = true;
                                                               e[++mlc] = \{u, head[v], 0, -cs\};
  for (int i = 0; i < n; ++i) {</pre>
                                                               head[v] = mlc;
  if (del[i] || v[i]) continue;
   g[i] += w[c][i];
                                                              void dfs(int u) {
                                                              vis[u] = 1;
                                                               for (int i = head[u], v; i; i = e[i].next)
if (!vis[v = e[i].to] and e[i].flow)
return make_pair(s, t);
                                                                 fa[v] = u, fae[v] = i, dfs(v);
int mincut(int n) {
int cut = 1e9;
                                                              ll phi(int x) {
memset(del, false, sizeof(del));
                                                               static ll pi[N];
for (int i = 0; i < n - 1; ++i) {</pre>
 int s, t; tie(s, t) = phase(n);
del[t] = true; cut = min(cut, g[t]);
                                                               if (x == -1) return 0;
                                                               if (vis[x] == visc) return pi[x];
                                                               return vis[x] = visc, pi[x] = phi(fa[x]) - e[fae[x]].
 for (int j = 0; j < n; ++j) {
                                                                  cost:
  w[s][j] += w[t][j]; w[j][s] += w[j][t];
 }
                                                              void pushflow(int x, ll &cost) {
}
                                                               int v = e[x ^ 1].to, u = e[x].to;
return cut;
                                                               ++visc;
                                                               while (v != -1) vis[v] = visc, v = fa[v];
4.7 GomoryHu Tree [7473bb]
                                                               while (u != -1 && vis[u] != visc)
                                                                vis[u] = visc, u = fa[u];
vector<tuple<int, int, int>> GomoryHu(int n){
                                                               vector<int> cyc;
vector<tuple<int, int, int>> rt;
                                                               int e2 = 0, pa = 2;
vector<int> g(n);
                                                               ll f = e[x].flow;
for (int i = 1; i < n; ++i) {</pre>
                                                               for (int i = e[x ^ 1].to; i != u; i = fa[i]) {
 int t = g[i];
                                                                cyc.push_back(fae[i]);
 flow.reset(); // clear flows on all edge
                                                                if (e[fae[i]].flow < f)</pre>
  rt.emplace_back(i, t, flow.max_flow(i, t));
                                                                 f = e[fae[e2 = i] ^ (pa = 0)].flow;
 flow.walk(i); // bfs points that connected to i (use
    edges with .cap > 0)
                                                               for (int i = e[x].to; i != u; i = fa[i]) {
  for (int j = i + 1; j < n; ++j)
                                                                cyc.push_back(fae[i] ^ 1);
  if (g[j]==t&&flow.connect(j)) // check if i can
                                                                if (e[fae[i] ^ 1].flow < f)</pre>
    reach j
                                                                 f = e[fae[e2 = i] ^ (pa = 1)].flow;
    g[j] = i;
                                                               cyc.push_back(x);
return rt;
                                                               for (int cyc_i : cyc) {
                                                                e[cyc_i].flow -= f, e[cyc_i ^ 1].flow += f;
4.8 Kuhn Munkres [2c09ed]
                                                                cost += 1ll * f * e[cyc_i].cost;
struct KM { // maximize, test @ UOJ 80
                                                               if (pa == 2) return;
int n, l, r; lld ans; // fl and fr are the match
                                                               int le = x ^ pa, l = e[le].to, o = e[le ^ 1].to;
vector<lld> hl, hr; vector<int> fl, fr, pre, q;
                                                               while (l != e2) {
void bfs(const auto &w, int s) {
                                                                vis[o] = 0;
 vector<int> vl(n), vr(n); vector<lld> slk(n, INF);
l = r = 0; vr[q[r++] = s] = true;
                                                                swap(le ^= 1, fae[o]), swap(l, fa[o]), swap(l, o);
  const auto check = [\&](int x) -> bool {
   if (vl[x] || slk[x] > 0) return true;
                                                              ll simplex() { // 1-based
   vl[x] = true; slk[x] = INF;
                                                               ll cost = 0;
   if (fl[x] != -1) return vr[q[r++] = fl[x]] = true;
                                                               memset(fa, -1, sizeof(fa)), dfs(1);
   while (x != -1) swap(x, fr[fl[x] = pre[x]]);
                                                               vis[1] = visc = 2, fa[1] = -1;
   return false;
                                                               for (int i = 2, pre = -1; i != pre; i = (i == mlc ? 2
                                                                  : i + 1))
  while (true) {
                                                                if (e[i].flow and e[i].cost < phi(e[i ^ 1].to) - phi(</pre>
   while (l < r)
                                                                  e[i].to))
    for (int x = 0, y = q[l++]; x < n; ++x) if (!vl[x])
```

pushflow(pre = i, cost);

return cost;

# 4.10 Minimum Cost Max Flow [6d1b01]

```
template <typename F, typename C> class MCMF {
static constexpr F INF_F = numeric_limits<F>::max();
static constexpr C INF_C = numeric_limits<C>::max();
struct E {
 int to, r;
 F f; C c;
 E() {}
 E(int a, int b, F x, C y)
  : to(a), r(b), f(x), c(y) {}
};
vector<vector<E>> g;
vector<pair<int, int>> f;
vector<bool> inq;
 vector<F> up; vector<C> d;
optional<pair<F, C>> step(int S, int T) {
  queue<int> q;
  for (q.push(S), d[S] = 0, up[S] = INF_F;
   not q.empty(); q.pop()) {
   int u = q.front(); inq[u] = false;
  if (up[u] == 0) continue;
for (int i = 0; i < int(g[u].size()); ++i) {</pre>
    auto e = g[u][i]; int v = e.to;
    if (e.f <= 0 or d[v] <= d[u] + e.c)
     continue:
    d[v] = d[u] + e.c; f[v] = {u, i};
    up[v] = min(up[u], e.f);
    if (not inq[v]) q.push(v);
    inq[v] = true;
  }
  if (d[T] == INF_C) return nullopt;
  for (int i = T; i != S; i = f[i].first) {
   auto &eg = g[f[i].first][f[i].second];
   eg.f -= up[T];
   g[eg.to][eg.r].f += up[T];
  return pair{up[T], d[T]};
public:
MCMF(int n) : g(n), f(n), inq(n), up(n), d(n, INF_C)  {}
void add_edge(int s, int t, F c, C w) {
 g[s].emplace_back(t, int(g[t].size()), c, w);
g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
pair<F, C> solve(int a, int b) {
 F c = 0; C w = 0;
 while (auto r = step(a, b)) {
   c += r->first, w += r->first * r->second;
   fill(inq.begin(), inq.end(), false);
   fill(d.begin(), d.end(), INF_C);
 return {c, w};
```

#### 4.11 Weighted Matching [94ca35]

```
#define pb emplace_back
#define rep(i, l, r) for (int i=(l); i<=(r); ++i)
struct WeightGraph { // 1-based
static const int inf = INT_MAX;
struct edge { int u, v, w; }; int n, nx;
vector<int> lab; vector<vector<edge>> g;
vector<int> slack, match, st, pa, S, vis;
vector<vector<int>> flo, flo_from; queue<int> q;
WeightGraph(int n_) : n(n_-), nx(n * 2), lab(nx + 1),
 g(nx + 1, vector < edge > (nx + 1)), slack(nx + 1)
  flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
 match = st = pa = S = vis = slack;
 rep(u, 1, n) rep(v, 1, n) g[u][v] = {u, v, 0};
int ED(edge e) {
 return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
 void update_slack(int u, int x, int &s) {
 if (!s || ED(g[u][x]) < ED(g[s][x])) s = u; }</pre>
void set_slack(int x) {
 slack[x] = 0;
  for (int u = 1; u <= n; ++u)
   if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
    update_slack(u, x, slack[x]);
void q_push(int x) {
```

```
if (x <= n) q.push(x);
else for (int y : flo[x]) q_push(y);
void set_st(int x, int b) {
st[x] = b;
if (x > n) for (int y : flo[x]) set_st(y, b);
vector<int> split_flo(auto &f, int xr) {
auto it = find(all(f), xr);
if (auto pr = it - f.begin(); pr % 2 == 1)
 reverse(1 + all(f)), it = f.end() - pr;
auto res = vector(f.begin(), it);
return f.erase(f.begin(), it), res;
void set_match(int u, int v) {
match[u] = g[u][v].v;
if (u <= n) return;</pre>
int xr = flo_from[u][g[u][v].u];
auto &f = flo[u], z = split_flo(f, xr);
rep(i, 0, int(z.size())-1) set_match(z[i], z[i ^ 1]);
set_match(xr, v); f.insert(f.end(), all(z));
void augment(int u, int v) {
for (;;) {
 int xnv = st[match[u]]; set_match(u, v);
 if (!xnv) return;
 set_match(xnv, st[pa[xnv]]);
 u = st[pa[xnv]], v = xnv;
}
int lca(int u, int v) {
static int t = 0; ++t;
for (++t; u || v; swap(u, v)) if (u) {
 if (vis[u] == t) return u;
 vis[u] = t; u = st[match[u]];
 if (u) u = st[pa[u]];
return 0;
void add_blossom(int u, int o, int v) {
int b = int(find(n + 1 + all(st), 0) - begin(st));
lab[b] = 0, S[b] = 0; match[b] = match[o];
vector<int> f = {0};
for (int x = u, y; x != o; x = st[pa[y]])
 f.pb(x), f.pb(y = st[match[x]]), q_push(y);
 reverse(1 + all(f));
for (int x = v, y; x != o; x = st[pa[y]])
 f.pb(x), f.pb(y = st[match[x]]), q_push(y);
 flo[b] = f; set_st(b, b);
for (int x = 1; x \le nx; ++x)
 g[b][x].w = g[x][b].w = 0;
 for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;</pre>
for (int xs : flo[b]) {
  for (int x = 1; x <= nx; ++x)
   if (g[b][x].w == 0 \mid \mid ED(g[xs][x]) < ED(g[b][x]))
   g[b][x] = g[xs][x], g[x][b] = g[x][xs];
 for (int x = 1; x \le n; ++x)
  if (flo_from[xs][x]) flo_from[b][x] = xs;
set_slack(b);
void expand_blossom(int b) {
for (int x : flo[b]) set_st(x, x);
int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
for (int x : split_flo(flo[b], xr)) {
 if (xs == -1) { xs = x; continue; }
 pa[xs] = g[x][xs].u; S[xs] = 1, S[x] = 0;
 slack[xs] = 0; set_slack(x); q_push(x); xs = -1;
for (int x : flo[b])
 if (x == xr) S[x] = 1, pa[x] = pa[b];
 else S[x] = -1, set_slack(x);
st[b] = 0;
bool on_found_edge(const edge &e) {
if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
 int nu = st[match[v]]; pa[v] = e.u; S[v] = 1;
 slack[v] = slack[nu] = 0; S[nu] = 0; q_push(nu);
} else if (S[v] == 0) {
 if (int o = lca(u, v)) add_blossom(u, o, v);
 else return augment(u, v), augment(v, u), true;
```

```
return false;
bool matching() {
  ranges::fill(S, -1); ranges::fill(slack, 0);
  q = queue<int>();
  for (int x = 1; x <= nx; ++x)</pre>
   if (st[x] == x && !match[x])
    pa[x] = 0, S[x] = 0, q_push(x);
  if (q.empty()) return false;
  for (;;) {
   while (q.size()) {
    int u = q.front(); q.pop();
    if (S[st[u]] == 1) continue;
    for (int v = 1; v <= n; ++v)</pre>
     if (g[u][v].w > 0 && st[u] != st[v]) {
      if (ED(g[u][v]) != 0)
       update_slack(u, st[v], slack[st[v]]);
      else if (on_found_edge(g[u][v])) return true;
     }
   int d = inf:
   for (int b = n + 1; b <= nx; ++b)
    if (st[b] == b && S[b] == 1)
    d = min(d, lab[b] / 2);
   for (int x = 1; x <= nx; ++x)
    if (int s = slack[x]; st[x] == x && s && S[x] <= 0)</pre>
     d = min(d, ED(g[s][x]) / (S[x] + 2));
   for (int u = 1; u <= n; ++u)</pre>
    if (S[st[u]] == 1) lab[u] += d;
    else if (S[st[u]] == 0) {
     if (lab[u] <= d) return false;</pre>
     lab[u] -= d;
   rep(b, n + 1, nx) if (st[b] == b && S[b] >= 0)
   lab[b] += d * (2 - 4 * \bar{S}[b]);
   for (int x = 1; x <= nx; ++x)</pre>
    if (int s = slack[x]; st[x] == x &&
      s \&\& st[s] != x \&\& ED(g[s][x]) == 0)
     if (on_found_edge(g[s][x])) return true;
   for (int b = n + 1; b <= nx; ++b)
    if (st[b] == b && S[b] == 1 && lab[b] == 0)
     expand_blossom(b);
  return false;
}
pair<lld, int> solve() {
 ranges::fill(match, 0);
  rep(u, 0, n) st[u] = u, flo[u].clear();
  int w_max = 0;
  rep(u, 1, n) rep(v, 1, n) {
  flo_from[u][v] = (u == v ? u : 0);
   w_max = max(w_max, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
 int n_matches = 0; lld tot_weight = 0;
while (matching()) ++n_matches;
  rep(u, 1, n) if (match[u] \&\& match[u] < u)
  tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
void set_edge(int u, int v, int w) {
  g[u][v].w = g[v][u].w = w; }
5
     Math
5.1 Common Bounds
```

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \backslash \{0\}} (-1)^{k+1} p(n-k(3k-1)/2) \approx 0.145/n \cdot \exp(2.56\sqrt{n})$$
 
$$\frac{n}{\max_{i \leq n} (d(i))} \frac{100 \text{ le3 le6 le9 le12 le15 le18}}{12 \text{ 32 240 l344 6720 26880 l03680}}$$

#### 5.2 Stirling Number

#### First Kind

 $S_1(n,k)$  counts the number of permutations of n elements with k disjoint

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$
$$x(x+1)\dots(x+n-1) = \sum_{k=0}^n S_1(n,k)x^k$$

$$g(x) = x(x+1)\dots(x+n-1) = \sum_{k=0}^{n} a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^{n} \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^{k} ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

#### Second Kind

 $S_2(n,k)$  counts the number of ways to partition a set of n elements into knonempty sets.  $S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$ 

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

# 5.3 ax+by=gcd [d0cbdd]

```
// ax+ny = 1, ax+ny == ax == 1 \pmod{n}
void exgcd(lld x, lld y, lld &g, lld &a, lld &b) {
 if (y == 0) g = x, a = 1, b = 0;
 else exgcd(y, x \% y, g, b, a), b = (x / y) * a;
```

# 5.4 Chinese Remainder [d69e74]

```
// please ensure r_i\in[0,m_i)
bool crt(lld &m1, lld &r1, lld m2, lld r2) {
  if (m2 > m1) swap(m1, m2), swap(r1, r2);
  lld g, a, b; exgcd(m1, m2, g, a, b);
  if ((r2 - r1) % g != 0) return false;
  m2 /= g; lld D = (r2 - r1) / g % m2 * a % m2;
r1 += (D < 0 ? D + m2 : D) * m1; m1 *= m2;
  assert (r1 >= 0 && r1 < m1);
  return true;
```

#### 5.5 DiscreteLog [86e463]

```
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
 // x^? \equiv y (mod M)
 Int t = 1, c = 0, g = 1;
for (Int M_ = M; M_ > 0; M_ >>= 1) g = g * x % M;
 for (g = gcd(g, M); t % g != 0; ++c) {
  if (t == y) return c;
  t = t * x % M;
 if (y % g != 0) return -1;
 t /= g, y /= g, M /= g;
 Int h = 0, gs = 1;

for (; h * h < M; ++h) gs = gs * x % M;
 unordered_map<Int, Int> bs;
 for (Int s = 0; s < h; bs[y] = ++s) y = y * x % M;
 for (Int s = 0; s < M; s += h) {</pre>
  t = t * gs % M;
  if (bs.count(t)) return c + s + h - bs[t];
 return -1;
}
```

# 5.6 Quadratic Residue [leabad]

```
int get_root(int n, int P) { // ensure 0 <= n < p</pre>
 if (P == 2 or n == 0) return n;
 auto check = [&](int x) {
  return modpow(x, (P - 1) / 2, P); };
if (check(n) != 1) return -1;
mt19937 rnd(7122); lld z = 1, w;
 while (check(w = (z * z - n + P) % P) != P - 1)
  z = rnd() \% P;
 const auto M = [P, w](auto &u, auto &v) {
  auto [a, b] = u; auto [c, d] = v;
  return make_pair((a * c + b * d % P * w) % P,
    (a * d + b * c) % P);
 pair<lld, lld> r(1, 0), e(z, 1);
 for (int w = (P + 1) / 2; w; w >>= 1, e = M(e, e))
  if (w & 1) r = M(r, e);
 return r.first; // sqrt(n) mod P where P is prime
```

#### Extended Euler

```
a^b \equiv \begin{cases} a^{(b \mod \varphi(m)) + \varphi(m)} & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases}
                                                                                                                                                                (\text{mod } m)
```

#### 5.8 Extended FloorSum

```
g(a,b,c,n) = \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor
                           \left( \left| \frac{a}{c} \right| \cdot \frac{n(n+1)(2n+1)}{6} + \left| \frac{b}{c} \right| \cdot \frac{n(n+1)}{2} \right)
                           +g(a \bmod c, b \bmod c, c, n),
                                                                                                       a \geq c \vee b \geq c
                                                                                                       n < 0 \lor a = 0
                           \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
                          -h(c, c-b-1, a, m-1)),
                                                                                                       otherwise
h(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^{2}
                          \left( \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right)
                            +\lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1)
                            +h(a \bmod c, b \bmod c, c, n)
                            +2\lfloor \tfrac{a}{c}\rfloor \cdot g(a \bmod c, b \bmod c, c, n)
                           +2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
                                                                                                       a \ge c \lor b \ge c
                                                                                                       n < 0 \lor a = 0
                            0,
                            nm(m+1) - 2g(c, c-b-1, a, m-1)
                           -2f(c, c-b-1, a, m-1) - f(a, b, c, n), otherwise
```

## 5.9 Extended Euclidean [e09892]

```
template <typename T>
auto euclid(lld a, lld b, lld c, lld n, T U, T R) {
b %= c;
if (a >= c)
 return euclid(a % c, b, c, n, U, mpow(U, a / c) * R);
lld m = (i128(a) * n + b) / c;
if (!m) return mpow(R, n);
return mpow(R, (c - b - 1) / a) * U *
       euclid(c, c - b - 1, a, m - 1, R, U) \star
       mpow(R, n - (i128(c) * m - b - 1) / a);
// time complexity is O(log max(a, c))
// 給定二維座標系上的一次函數 $y = (ax + b) / c$
// 維護一個矩陣 $A = I$, 考慮 $x \in [0, n)$
// 每次向右穿過網格的垂直線時,乘上一個矩陣 $R$
// 每次向上穿過網格的水平線時,乘上一個矩陣 $U$
// 若剛好經過一個整點,那麼先乘 $U$ 再乘 $R$
```

#### 5.10 FloorSum ffb59171

```
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
// @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
llu ans = 0;
while (true) {
 if (a >= m) ans += n*(n-1)/2 * (a/m), a %= m;
  if (b >= m) ans += n * (b/m), b %= m;
  if (llu y_max = a * n + b; y_max >= m) {
  n = (llu)(y_max / m), b = (llu)(y_max % m);
   swap(m, a);
 } else break;
}
return ans;
lld floor_sum(lld n, lld m, lld a, lld b) {
llu ans = 0:
if (a < 0) {
 llu a2 = (a \% m + m), d = (a2 - a) / m;
 ans -= 1ULL * n * (n - 1) / 2 * d; a = a2;
if (b < 0) {
 llu b2 = (b \% m + m), d = (b2 - b) / m;
 ans -= 1ULL * n * d; b = b2;
return ans + floor_sum_unsigned(n, m, a, b);
```

# 5.11 ModMin [253e4d]

```
// min{k | l <= ((ak) mod m) <= r}
optional<llu> mod_min(u32 a, u32 m, u32 l, u32 r) {
   if (a == 0) return l ? nullopt : 0;
   if (auto k = llu(l + a - 1) / a; k * a <= r)
      return k;
   auto b = m / a, c = m % a;
   if (auto y = mod_min(c, a, a - r % a, a - l % a))
    return (l + *y * c + a - 1) / a + *y * b;
   return nullopt;
}</pre>
```

### **5.12** FWT [c5167a]

```
/* or convolution:
    x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
    * and convolution:
    x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
void fwt(int x[], int N, bool inv = false) {
    for (int d = 1; d < N; d <<= 1)
        for (int s = 0; s < N; s += d * 2)
        for (int i = s; i < s + d; i++) {
            int j = i + d, ta = x[i], tb = x[j];
            x[i] = modadd(ta, tb);
            x[j] = modsub(ta, tb);
        }
    if (inv) {
        const int invn = modinv(N);
        for (int i = 0; i < N; i++)
            x[i] = modmul(x[i], invn);
    }
}</pre>
```

### 5.13 Packed FFT [321552]

```
int round2k(size_t n) {
 int sz = 1; while (sz < int(n)) sz *= 2; return sz; }</pre>
VL convolution(const VI &a, const VI &b) {
  const int sz = round2k(a.size() + b.size() - 1);
 // Should be able to handle N <= 10^5, C <= 10^4
 vector<P> v(sz);
 for (size_t i = 0; i < a.size(); i++) v[i].RE(a[i]);</pre>
 for (size_t i = 0; i < b.size(); i++) v[i].IM(b[i]);</pre>
 fft(v.data(), sz, /*inv=*/false);
 auto rev = v; reverse(1 + all(rev));
 for (int i = 0; i < sz; i++) {</pre>
  P A = (v[i] + conj(rev[i])) / P(2, 0);
  P B = (v[i] - conj(rev[i])) / P(0, 2);
  v[i] = A * B;
 VL c(sz); fft(v.data(), sz, /*inv=*/true);
for (int i = 0; i < sz; ++i) c[i] = roundl(RE(v[i]));</pre>
 return c;
VI convolution_mod(const VI &a, const VI &b) {
 const int sz = round2k(a.size() + b.size() - 1);
 vector<P> fa(sz), fb(sz);
 for (size_t i = 0; i < a.size(); ++i)</pre>
  fa[i] = P(a[i] & ((1 << 15) - 1), a[i] >> 15);
 for (size_t i = 0; i < b.size(); ++i)</pre>
  fb[i] = P(b[i] & ((1 << 15) - 1), b[i] >> 15);
 fft(fa.data(), sz); fft(fb.data(), sz);
 auto rfa = fa; reverse(1 + all(rfa));
 for (int i = 0; i < sz; ++i) fa[i] *= fb[i];</pre>
 for (int i = 0; i < sz; ++i) fb[i] *= conj(rfa[i]);</pre>
 fft(fa.data(), sz, true); fft(fb.data(), sz, true);
 vector<int> res(sz);
 for (int i = 0; i < sz; ++i) {</pre>
  lld A = (lld)roundl(RE((fa[i] + fb[i]) / P(2, 0)));
  lld C = (lld)roundl(IM((fa[i] - fb[i]) / P(0, 2)));
lld B = (lld)roundl(IM(fa[i])); B %= p; C %= p;
  res[i] = (A + (B << 15) + (C << 30)) % p;
 return res;
} // test @ yosupo judge with long double
5.14 CRT for arbitrary mod [e4dde7]
```

```
const int mod = 10000000007;
const int M1 = 985661441; // G = 3 for M1, M2, M3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(lld A, lld B, lld C) {
   static_assert (M1 < M2 && M2 < M3);
   constexpr lld r12 = modpow(M1, M2-2, M2);
   constexpr lld r13 = modpow(M1, M3-2, M3);
   constexpr lld r23 = modpow(M2, M3-2, M3);
   constexpr lld M1M2 = 1LL * M1 * M2 % mod;
   B = (B - A + M2) * r12 % M2;
   C = (C - A + M3) * r13 % M3;
   C = (C - B + M3) * r23 % M3;
   return (A + B * M1 + C * M1M2) % mod;
}</pre>
```

# **5.15** NTT / FFT [03190d]

```
template <int mod, int G, int maxn> struct NTT {
  static_assert (maxn == (maxn & -maxn));
```

Poly Pow(Poly a, lld M) { // period mod\*(mod-1)

```
assert(!a.empty() && a[0] != 0);
 int roots[maxn];
                                                                 const int N = int(a.size()); // mod x^N
 NTT () {
  int r = modpow(G, (mod - 1) / maxn);
                                                                 const auto imul = [&a](int s) {
  for (int i = maxn >> 1; i; i >>= 1) {
                                                                  for (int &x: a) x = modmul(x, s); }; int c = a[0];
   roots[i] = 1;
                                                                 imul(modinv(c)); a = Ln(a); imul(M % mod);
   for (int j = 1; j < i; j++)</pre>
                                                                 a = Exp(a); imul(modpow(c, M % (mod - 1)));
    roots[i + j] = modmul(roots[i + j - 1], r);
                                                                 return a;
   r = modmul(r, r);
// for (int j = 0; j < i; j++) // FFT (tested)</pre>
                                                                Poly Sqrt(const Poly &v) { // need: QuadraticResidue assert(!v.empty() && v[0] != 0);
       roots[i+j] = polar < llf > (1, PI * j / i);
                                                                 const int r = get_root(v[0]); assert(r != -1);
                                                                 return Newton(v, [r](int x) { return r; },
 // n must be 2^k, and 0 <= F[i] < mod
void operator()(int F[], int n, bool inv = false) {</pre>
                                                                  [](Poly &X, Poly &A, int sz) {
                                                                   auto Y = X; Y.resize(sz / 2);
                                                                   auto B = Mul(A, Inv(Y), sz);
  for (int i = 0, j = 0; i < n; i++) {
   if (i < j) swap(F[i], F[j]);
for (int k = n>>1; (j^=k) < k; k>>=1);
                                                                   for (int i = 0, inv2 = mod / 2 + 1; i < sz; i++)</pre>
                                                                    X[i] = modmul(inv2, modadd(X[i], B[i])); });
  for (int s = 1; s < n; s *= 2) {
                                                                Poly Mul(auto &&a, auto &&b) {
   for (int i = 0; i < n; i += s * 2) {</pre>
                                                                 const int n = a.size() + b.size() - 1;
                                                                 auto R = Mul(a, b, bit_ceil(n));
    for (int j = 0; j < s; j++) {</pre>
     int a = F[i+j], b = modmul(F[i+j+s], roots[s+j]);
                                                                 return R.resize(n), R;
     F[i+j] = modadd(a, b); // a + b
     F[i+j+s] = modsub(a, b); // a - b
                                                                Poly MulT(Poly a, Poly b, int k) {
                                                                 assert(b.size()); reverse(all(b)); auto R = Mul(a, b);
   }
                                                                 R = vector(R.begin() + b.size() - 1, R.end());
                                                                 return R.resize(k), R;
  if (inv) {
   int iv = modinv(n);
                                                                Poly Eval(const Poly &f, const Poly &x) {
   for (int i = 0; i < n; i++) F[i] = modmul(F[i], iv);</pre>
                                                                 if (f.empty()) return vector(x.size(), 0);
   reverse(F + 1, F + n);
                                                                 const int n = int(max(x.size(), f.size()));
                                                                 auto q = vector(n * 2, Poly(2, 1)); Poly ans(n);
 }
                                                                 fi(0, x.size()) q[i + n][1] = modsub(0, x[i]);
                                                                 for (int i = n - 1; i > 0; i--)
};
                                                                  q[i] = Mul(q[i << 1], q[i << 1 | 1]);
      Formal Power Series [2bc0d3]
                                                                 q[1] = MulT(f, Inv(q[1]), n);
                                                                 for (int i = 1; i < n; i++) {
#define fi(l, r) for (size_t i = (l); i < (r); ++i)
using Poly = vector<int>;
                                                                  auto L = q[i << 1], R = q[i << 1 | 1];</pre>
                                                                  q[i << 1 | 0] = MulT(q[i], R, L.size());
q[i << 1 | 1] = MulT(q[i], L, R.size());</pre>
auto Mul(auto a, auto b, size_t sz) {
 a.resize(sz), b.resize(sz);
 ntt(a.data(), sz); ntt(b.data(), sz);
 fi(0, sz) a[i] = modmul(a[i], b[i]);
                                                                 for (int i = 0; i < n; i++) ans[i] = q[i + n][0];</pre>
 return ntt(a.data(), sz, true), a;
                                                                 return ans.resize(x.size()), ans;
Poly Newton(const Poly &v, auto &&init, auto &&iter) { Poly Q = \{ init(v[0]) \}; 
                                                                pair<Poly, Poly> DivMod(const Poly &A, const Poly &B) {
                                                                 assert(!B.empty() && B.back() != 0);
                                                                 if (A.size() < B.size()) return {{}}, A};</pre>
 for (int sz = 2; Q.size() < v.size(); sz *= 2) {</pre>
                                                                 const int sz = A.size() - B.size() + 1;
 Poly A{begin(v), begin(v) + min(sz, int(v.size()))};
  A.resize(sz * 2), Q.resize(sz * 2);
                                                                 Poly X = B; reverse(all(X)); X.resize(sz);
  iter(Q, A, sz * 2); Q.resize(sz);
                                                                 Poly Y = A; reverse(all(Y)); Y.resize(sz);
                                                                 Poly Q = Mul(Inv(X), Y);
 return Q.resize(v.size()), Q;
                                                                 Q.resize(sz); reverse(all(Q)); X = Mul(Q, B); Y = A;
                                                                 fi(0, Y.size()) Y[i] = modsub(Y[i], X[i]);
Poly Inv(const Poly &v) { // v[0] != 0
                                                                 while (Y.size() && Y.back() == 0) Y.pop_back();
                                                                 while (Q.size() && Q.back() == 0) Q.pop_back();
 return Newton(v, modinv,
  [](Poly &X, Poly &A, int sz) {
                                                                 return {Q, Y};
   ntt(X.data(), sz), ntt(A.data(), sz);
                                                                 // empty means zero polynomial
                                                                int LinearRecursionKth(Poly a, Poly c, int64_t k) {
   for (int i = 0; i < sz; i++)</pre>
    X[i] = modmul(X[i], modsub(2, modmul(X[i], A[i])));
                                                                 const auto d = a.size(); assert(c.size() == d + 1);
                                                                 const int sz = bit_ceil(2 * d + 1), o = sz / 2;
   ntt(X.data(), sz, true); });
                                                                 Poly q = c; for (int &x: q) x = modsub(0, x); q[0]=1;
                                                                 Poly p = Mul(a, q); p.resize(sz); q.resize(sz);
Poly Dx(Poly A) {
 fi(1, A.size()) A[i - 1] = modmul(i, A[i]);
                                                                 for (int r; r = (k & 1), k; k >>= 1) {
                                                                  fill(d + all(p), 0); fill(d + 1 + all(q), 0);
 return A.empty() ? A : (A.pop_back(), A);
                                                                  ntt(p.data(), sz); ntt(q.data(), sz);
Poly Sx(Poly A) {
                                                                  for(int i = 0; i < sz; i++)</pre>
 A.insert(A.begin(), 0);
                                                                   p[i] = modmul(p[i], q[(i + o) & (sz - 1)]);
 fi(1, A.size()) A[i] = modmul(modinv(i), A[i]);
                                                                  for(int i = 0, j = 0; j < sz; i++, j++)</pre>
                                                                   q[i] = q[j] = modmul(q[i], q[j]);
 return A;
                                                                  ntt(p.data(), sz, true); ntt(q.data(), sz, true);
for (int i = 0; i < d; i++) p[i] = p[i << 1 | r];</pre>
Poly Ln(const Poly &A) { // coef[0] == 1; res[0] == 0
 auto B = Sx(Mul(Dx(A), Inv(A), bit_ceil(A.size()*2)));
                                                                  for (int i = 0; i <= d; i++) q[i] = q[i << 1];
                                                                 } // Bostan-Mori
 return B.resize(A.size()), B;
                                                               return modmul(p[0], modinv(q[0]));
} // a_n = \sum_j a_n(n-j), c_0 \text{ is not important}
Poly Exp(const Poly &v) { // coef[0] == 0; res[0] == 1
 return Newton(v, [](int x) { return 1; },
                                                                5.17
                                                                       Partition Number [9bb845]
  [](Poly &X, Poly &A, int sz) {
   auto Y = X; Y.resize(sz / 2); Y = Ln(Y);
                                                                ans[0] = tmp[0] = 1;
   fi(0, Y.size()) Y[i] = modsub(A[i], Y[i]);
                                                                for (int i = 1; i * i <= n; i++) {</pre>
                                                                 for (int rep = 0; rep < 2; rep++)
for (int j = i; j <= n - i * i; j++)
modadd(tmp[j], tmp[j-i]);</pre>
   Y[0] = modadd(Y[0], 1); X = Mul(X, Y, sz); );
```

```
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 for (int j = i * i; j <= n; j++)</pre>
 modadd(ans[j], tmp[j - i * i]);
5.18 Pi Count (+Linear Sieve) [8a4382]
static constexpr int N = 1000000 + 5;
lld pi[N]; vector<int> primes; bool sieved[N];
lld cube_root(lld x) {
lld s = cbrt(x - 0.1L);
while (s * s * s <= x) ++s;
return s - 1;
lld square_root(lld x) {
lld s = sqrt(x - 0.1L);
while (s * s <= x) ++s;
return s - 1;
void init() {
primes.reserve(N);
for (int i = 2; i < N; i++) {
 if (!sieved[i]) primes.push_back(i);
 pi[i] = !sieved[i] + pi[i - 1];
  for (int p : primes) {
  if (i * p >= N) break;
   sieved[p * i] = true;
   if (i % p == 0) break;
 }
primes.insert(primes.begin(), 1);
lld phi(lld m, lld n) {
static constexpr int MM = 80000, NN = 500;
 static lld val[MM][NN];
if (m<MM && n<NN && val[m][n]) return val[m][n] - 1;</pre>
if (n == 0) return m;
 if (primes[n] >= m) return 1;
lld ret = phi(m, n - 1) - phi(m / primes[n], n - 1);
if (m < MM && n < NN) val[m][n] = ret + 1;</pre>
return ret;
lld pi_count(lld);
lld P2(lld m, lld n) {
lld sm = square_root(m), ret = 0;
for (lld i = n + 1; primes[i] <= sm; i++)</pre>
 ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
return ret;
lld pi_count(lld m) {
if (m < N) return pi[m];</pre>
lld n = pi_count(cube_root(m));
return phi(m, n) + n - 1 - P2(m, n);
5.19 Miller Rabin [fbd812]
bool isprime(llu x) {
auto witn = [&](llu a, int t) {
  for (llu a2; t--; a = a2) {
   a2 = mmul(a, a, x);
   if (a2 == 1 && a != 1 && a != x - 1) return true;
 }
 return a != 1;
 if (x <= 2 || ~x & 1) return x == 2;</pre>
 int t = countr_zero(x-1); llu odd = (x-1) >> t;
 for (llu m:
  {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
  if (m % x != 0 && witn(mpow(m % x, odd, x), t))
   return false:
return true:
} // test @ luogu 143 & yosupo judge
5.20 Pollard Rho [57ad88]
// does not work when n is prime or n == 1
// return any non-trivial factor
llu pollard_rho(llu n) {
static mt19937_64 rnd(120821011);
 if (!(n & 1)) return 2;
llu y = 2, z = y, c = rnd() % n, p = 1, i = 0, t;
auto f = [&](llu x) {
  return madd(mmul(x, x, n), c, n); };
```

p = mmul(msub(z = f(f(z)), y = f(y), n), p, n);

**do** {

```
if (++i &= 63) if (i == (i & -i)) t = gcd(p, n);
 } while (t == 1);
 return t == n ? pollard_rho(n) : t;
} // test @ yosupo judge
5.21 Berlekamp Massey [a94d00]
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
 vector<T> d(output.size() + 1), me, he;
 for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
  for (size_t j = 0; j < me.size(); ++j)
d[i] += output[i - j - 2] * me[j];</pre>
  if ((d[i] -= output[i - 1]) == 0) continue;
  if (me.empty()) {
   me.resize(f = i);
   continue;
  }
  vector<T> o(i - f - 1);
  T k = -d[i] / d[f]; o.push_back(-k);
  for (T x : he) o.push_back(x * k);
  if (o.size() < me.size()) o.resize(me.size());</pre>
  for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
  if (i-f+he.size() >= me.size()) he = me, f = i;
  me = o;
 return me;
5.22 Gauss Elimination [c131f4]
void gauss(vector<vector<llf>> &d) {
 int n = d.size(), m = d[0].size();
for (int i = 0; i < m; ++i) {</pre>
  int p = -1;
  for (int j = i; j < n; ++j)
if (abs(d[j][i]) > eps)
    if (p == -1 || abs(d[j][i]) > abs(d[p][i]))
  p = j;
if (p == -1) continue;
  for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);</pre>
  for (int j = 0; j < n; ++j) if (j != i) {</pre>
   llf z = d[j][i] / d[i][i];
   for (int k = 0; k < m; ++k)
    d[j][k] = z * d[i][k];
5.23
       Charateristic Polynomial [ff2159]
#define rep(x, y, z) for (int x=y; x<z; x++)</pre>
using VI = vector<int>; using VVI = vector<VI>;
void Hessenberg(VVI &H, int N) {
 for (int i = 0; i < N - 2; ++i) {
  for (int j = i + 1; j < N; ++j) if (H[j][i]) {</pre>
   rep(k, i, N) swap(H[i+1][k], H[j][k]);
   rep(k, 0, N) swap(H[k][i+1], H[k][j]);
   break:
  if (!H[i + 1][i]) continue;
  for (int j = i + 2; j < N; ++j) {
   int co = mul(modinv(H[i + 1][i]), H[j][i]);
   rep(k, i, N) subeq(H[j][k], mul(H[i+1][k], co));
   rep(k, 0, N) addeq(H[k][i+1], mul(H[k][j], co));
 }
VI CharacteristicPoly(VVI &A) {
 int N = (int)A.size(); Hessenberg(A, N);
 VVI P(N + 1, VI(N + 1)); P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
  rep(j, 0, i+1) P[i][j] = j ? P[i-1][j-1] : 0;</pre>
  for (int j = i - 1, val = 1; j >= 0; --j) {
  int co = mul(val, A[j][i - 1]);
   rep(k, 0, j+1) subeq(P[i][k], mul(P[j][k], co));
   if (j) val = mul(val, A[j][j - 1]);
 if (N & 1) for (int &x: P[N]) x = sub(0, x);
 return P[N]; // test: 2021 PTZ Korea K
```

5.24 Simplex [c9c93b]

```
namespace simplex {
                                                               llf F(llf L, llf R, llf v, llf eps) {
                                                                llf M = (L + R) / 2, vl = simp(L, M), vr = simp(M, R);
// maximize c^Tx under Ax \le B and x \ge 0
// return VD(n, -inf) if the solution doesn't exist
                                                                if (abs(vl + vr - v) <= 15 * eps)
// return VD(n, +inf) if the solution is unbounded
                                                                 return vl + vr + (vl + vr - v) / 15.0;
using VD = vector<llf>;
                                                                return F(L, M, vl, eps / 2.0) +
using VVD = vector<vector<llf>>;
                                                                     F(M, R, vr, eps / 2.0);
const llf eps = 1e-9, inf = 1e+9;
                                                               } // call F(l, r, simp(l, r), 1e-6)
int n, m; VVD d; vector<int> p, q;
                                                                     Geometry
void pivot(int r, int s) {
  llf inv = 1.0 / d[r][s];
                                                               6.1 Basic Geometry [e4a147]
for (int i = 0; i < m + 2; ++i)
 for (int j = 0; j < n + 2; ++j)
                                                               #define IM imag
   if (i != r && j != s)
                                                               #define RE real
    d[i][j] -= d[r][j] * d[i][s] * inv;
                                                               using lld = int64_t;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;</pre>
                                                               using llf = long double;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
                                                               using PT = std::complex<lld>;
d[r][s] = inv; swap(p[r], q[s]);
                                                               using PTF = std::complex<llf>;
                                                               using P = PT;
bool phase(int z) {
                                                               llf abs(P p) { return sqrtl(norm(p)); }
 int x = m + z;
                                                               PTF toPTF(PT p) { return PTF{RE(p), IM(p)}; }
while (true) {
                                                               int sgn(lld x) \{ return (x > 0) - (x < 0); \}
  int s = -1;
                                                               lld dot(P a, P b) { return RE(conj(a) * b); }
  for (int i = 0; i <= n; ++i) {</pre>
                                                               lld cross(P a, P b) { return IM(conj(a) * b); }
  if (!z && q[i] == -1) continue;
                                                               int ori(P a, P b, P c) {
   if (s == -1 || d[x][i] < d[x][s]) s = i;
                                                                return sgn(cross(b - a, c - a));
  if (s == -1 || d[x][s] > -eps) return true;
                                                               int quad(P p) {
 int r = -1;
for (int i = 0; i < m; ++i) {</pre>
                                                                return (IM(p) == 0) // use sgn for PTF
                                                                 ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
   if (d[i][s] < eps) continue;</pre>
   if (r == -1 ||
                                                               int argCmp(P a, P b) {
   d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
                                                                // returns 0/+-1, starts from theta = -PI
                                                                int qa = quad(a), qb = quad(b);
                                                                if (qa != qb) return sgn(qa - qb);
  if (r == -1) return false;
  pivot(r, s);
                                                                return sgn(cross(b, a));
                                                               P rot90(P p) { return P{-IM(p), RE(p)}; }
VD solve(const VVD &a, const VD &b, const VD &c) {
                                                               template <typename V> llf area(const V & pt) {
                                                                lld ret = 0;
m = (int)b.size(), n = (int)c.size();
d = VVD(m + 2, VD(n + 2));
                                                                for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
for (int i = 0; i < m; ++i)</pre>
                                                                 ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
                                                                return ret / 2.0;
p.resize(m), q.resize(n + 1);
for (int i = 0; i < m; ++i)</pre>
                                                               template <typename V> PTF center(const V & pt) {
 p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
                                                                P ret = 0; lld A = 0;
 for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];</pre>
                                                                for (int i = 1; i + 1 < (int)pt.size(); i++) {</pre>
q[n] = -1, d[m + 1][n] = 1;
                                                                 lld cur = cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 int r = 0;
                                                                 ret += (pt[i] + pt[i + 1] + pt[0]) * cur; A += cur;
 for (int i = 1; i < m; ++i)</pre>
  if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
                                                                return toPTF(ret) / llf(A * 3);
if (d[r][n + 1] < -eps) {
 pivot(r, n);
                                                               PTF project(PTF p, PTF q) { // p onto q
 if (!phase(1) || d[m + 1][n + 1] < -eps)</pre>
                                                                return dot(p, q) * q / dot(q, q); // dot<llf>
  return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
  int s = min_element(d[i].begin(), d[i].end() - 1)
                                                               6.2 2D Convex Hull [ecba37]
        - d[i].begin();
                                                               // from NaCl, counterclockwise, be careful of n<=2
  pivot(i, s);
                                                               vector<P> convex_hull(vector<P> v) {
 }
                                                                sort(all(v)); // by X then Y
                                                                 if (v[0] == v.back()) return {v[0]};
if (!phase(0)) return VD(n, inf);
                                                                int t = 0, s = 1; vector<P> h(v.size() + 1);
VD x(n);
                                                                for (int _ = 2; _--; s = t--, reverse(all(v)))
for (int i = 0; i < m; ++i)</pre>
                                                                 for (P p : v) {
 if (p[i] < n) x[p[i]] = d[i][n + 1];</pre>
                                                                  while (t>s && ori(p, h[t-1], h[t-2]) >= 0) t--;
 return x;
                                                                  h[t++] = p;
5.25 Simplex Construction
                                                                return h.resize(t), h;
Standard form: maximize \sum_{1 < i < n} c_i x_i such that for all 1 \le j \le m,
\sum_{1 \le i \le n} A_{ji} x_i \le b_j and x_i \ge 0 for all 1 \le i \le n.
  \bar{\mathsf{l}}. In case of minimization, let c_i' = -c_i
                                                               6.3 2D Farthest Pair [8b5844]
 2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} A_{ji} x_i \leq -b_j
3. \sum_{1 \leq i \leq n} A_{ji} x_i = b_j
                                                                // p is CCW convex hull w/o colinear points
                                                               int n = (int)p.size(), pos = 1; lld ans = 0;
       • \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
                                                               for (int i = 0; i < n; i++) {</pre>
       • \sum_{1 \le i \le n} A_{ji} x_i \ge b_j
                                                                P = p[(i + 1) \% n] - p[i];
                                                                while (cross(e, p[(pos + 1) % n] - p[i]) >
    If x_i has no lower bound, replace x_i with x_i - x_i'
5.26 Adaptive Simpson [09669e]
                                                                     cross(e, p[pos] - p[i]))
                                                                 pos = (pos + 1) % n;
llf simp(llf l, llf r) {
                                                                for (int j: {i, (i + 1) % n})
llf m = (l + r) / 2;
                                                                 ans = max(ans, norm(p[pos] - p[j]));
return (f(l) + f(r) + 4.0 * f(m)) * (r - l) / 6.0;
```

} // tested @ AOJ CGL\_4\_B

### 6.4 MinMax Enclosing Rect [e4470c]

```
// from 8BQube, plz ensure p is strict convex hull
const llf INF = 1e18, qi = acos(-1) / 2 * 3;
pair<llf, llf> solve(const vector<P> &p) {
llf mx = 0, mn = INF; int n = (int)p.size();
for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {</pre>
#define Z(v) (p[(v) % n] - p[i])
  P = Z(i + 1);
  while (cross(e, Z(u + 1)) > cross(e, Z(u))) ++u;
  while (dot(e, Z(r + 1)) > dot(e, Z(r))) ++r;
  if (!i) l = r + 1;
  while (dot(e, Z(l + 1)) < dot(e, Z(l))) ++l;</pre>
  P D = p[r \% n] - p[l \% n];
  llf H = cross(e, Z(u)) / llf(norm(e));
  mn = min(mn, dot(e, D) * H);
  llf B = sqrt(norm(D)) * sqrt(norm(Z(u)));
  llf deg = (qi - acos(dot(D, Z(u)) / B)) / 2;
 mx = max(mx, B * sin(deg) * sin(deg));
 return {mn, mx};
} // test @ UVA 819
```

# 6.5 Minkowski Sum [602806]

```
// A, B are strict convex hull rotate to min by (X, Y)
vector<P> Minkowski(vector<P> A, vector<P> B) {
 const int N = (int)A.size(), M = (int)B.size();
 vector<P> sa(N), sb(M), C(N + M + 1);
 for (int i = 0; i < N; i++) sa[i] = A[(i+1)%N]-A[i];
 for (int i = 0; i < M; i++) sb[i] = B[(i+1)%M]-B[i];</pre>
 C[0] = A[0] + B[0];
 for (int i = 0, j = 0; i < N || j < M; ) {
  P e = (j>=M || (i<N && cross(sa[i], sb[j])>=0))
   ? sa[i++] : sb[j++];
  C[i + j] = e;
 partial_sum(all(C), C.begin()); C.pop_back();
 return convex_hull(C); // just to remove colinear
```

# 6.6 Segment Intersection [60d016]

```
struct Seg { // closed segment
P st, dir; // represent st + t*dir for 0 \le t \le 1
Seg(P s, P e) : st(s), dir(e - s) {}
static bool valid(lld p, lld q) {
  // is there t s.t. 0 <= t <= 1 && qt == p ?
  if (q < 0) q = -q, p = -p;
  return 0 <= p && p <= q;
vector<P> ends() const { return { st, st + dir }; }
template <typename T> bool isInter(T A, P p) {
if (A.dir == P(0)) return p == A.st; // BE CAREFUL
return cross(p - A.st, A.dir) == 0 &&
 T::valid(dot(p - A.st, A.dir), norm(A.dir));
template <typename U, typename V>
bool isInter(U A, V B) {
if (cross(A.dir, B.dir) == 0) { // BE CAREFUL
 for (P p: A.ends()) res |= isInter(B, p);
 for (P p: B.ends()) res |= isInter(A, p);
  return res:
P D = B.st - A.st; lld C = cross(A.dir, B.dir);
return U::valid(cross(D, B.dir), C) &&
  V::valid(cross(D, A.dir), C);
```

#### 6.7 Half Plane Intersection [45e909]

```
struct Line {
P st, ed, dir;
Line (P s, P e) : st(s), ed(e), dir(e - s) {}
}; using LN = const Line &;
PTF intersect(LN A, LN B) {
llf t = cross(B.st - A.st, B.dir) /
  llf(cross(A.dir, B.dir));
return toPTF(A.st) + toPTF(A.dir) * t; // C^3 / C^2
bool cov(LN l, LN A, LN B) {
i128 u = cross(B.st-A.st, B.dir);
i128 v = cross(A.dir, B.dir);
// ori(l.st, l.ed, A.st + A.dir*(u/v)) <= 0?
```

```
i128 x = RE(A.dir) * u + RE(A.st - l.st) * v;
 i128 y = IM(A.dir) * u + IM(A.st - l.st) * v;
 return sgn(x*IM(l.dir) - y*RE(l.dir)) * sgn(v) >= 0;
} // x, y are C^3, also sgn<i128> is needed
bool operator<(LN a, LN b) {</pre>
 if (int c = argCmp(a.dir, b.dir)) return c == -1;
 return ori(a.st, a.ed, b.st) < 0;</pre>
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
// the half plane is the LHS when going from st to ed
llf HPI(vector<Line> &q) {
 sort(q.begin(), q.end());
 int n = (int)q.size(), l = 0, r = -1;
for (int i = 0; i < n; i++) {</pre>
  if (i && !argCmp(q[i].dir, q[i-1].dir)) continue;
  while (l < r && cov(q[i], q[r-1], q[r])) --r;</pre>
  while (l < r && cov(q[i], q[l], q[l+1])) ++l;</pre>
  q[++r] = q[i];
 while (l < r && cov(q[l], q[r-1], q[r])) --r;</pre>
 while (l < r && cov(q[r], q[l], q[l+1])) ++l;</pre>
 n = r - l + 1; // q[l .. r] are the lines
 if (n <= 1 || !argCmp(q[l].dir, q[r].dir)) return 0;</pre>
 vector<PTF> pt(n);
 for (int i = 0; i < n; i++)</pre>
  pt[i] = intersect(q[i+l], q[(i+1)%n+l]);
 return area(pt);
} // test @ 2020 Nordic NCPC : BigBrother
```

# 6.8 SegmentDist (Sausage) [9d8603]

```
'/ be careful of abs<complex<int>> (replace _abs below)
llf PointSegDist(P A, Seg B) {
 if (B.dir == P(0)) return _abs(A - B.st);
 if (sgn(dot(A - B.st, B.dir)) *
   sgn(dot(A - B.ed, B.dir)) <= 0)</pre>
  return abs(cross(A - B.st, B.dir)) / _abs(B.dir);
 return min(_abs(A - B.st), _abs(A - B.ed));
llf SegSegDist(const Seg &s1, const Seg &s2) {
 if (isInter(s1, s2)) return 0;
 return min({
   PointSegDist(s1.st, s2),
   PointSegDist(s1.ed, s2),
   PointSegDist(s2.st, s1),
   PointSegDist(s2.ed, s1) });
} // test @ QOJ2444 / PTZ19 Summer.D3
```

#### 6.9 Rotating Sweep Line [8aff27]

```
struct Event {
 Pd; int u, v;
 bool operator<(const Event &b) const {</pre>
  return sgn(cross(d, b.d)) > 0; }
P makePositive(P z) { return cmpxy(z, 0) ? -z : z; }
void rotatingSweepLine(const vector<P> &p) {
 const int n = int(p.size());
 vector<Event> e; e.reserve(n * (n - 1) / 2);
 for (int i = 0; i < n; i++)</pre>
  for (int j = i + 1; j < n; j++)</pre>
   e.emplace_back(makePositive(p[i] - p[j]), i, j);
 sort(all(e));
 vector<int> ord(n), pos(n);
 iota(all(ord), 0);
 sort(all(ord), [&p](int i, int j) {
  return cmpxy(p[i], p[j]); });
 for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
 const auto makeReverse = [](auto &v) {
  sort(all(v)); v.erase(unique(all(v)), v.end());
  vector<pair<int,int>> segs;
  for (size_t i = 0, j = 0; i < v.size(); i = j) {</pre>
   for (; j < v.size() && v[j] - v[i] <= j - i; j++);
segs.emplace_back(v[i], v[j - 1] + 1 + 1);</pre>
  return segs;
 for (size_t i = 0, j = 0; i < e.size(); i = j) {</pre>
  /* do here */
  vector<size_t> tmp;
  for (; j < e.size() && !(e[i] < e[j]); j++)</pre>
   tmp.push_back(min(pos[e[j].u], pos[e[j].v]));
  for (auto [l, r] : makeReverse(tmp)) {
```

reverse(ord.begin() + l, ord.begin() + r);

```
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   for (int t = l; t < r; t++) pos[ord[t]] = t;</pre>
                                                            vector<PTF> intersectPoint(Cir a, Cir b) {
                                                             llf d = abs(a.o - b.o);
  }
 }
}
      Polygon Cut [e9bdd1]
6.10
                                                             PTF u = dir * d1 + b.o;
using P = PTF;
vector<P> cut(const vector<P>& poly, P s, P e) {
                                                             return {u + v, u - v};
 vector<P> res;
                                                            } // test @ AOJ CGL probs
 for (size_t i = 0; i < poly.size(); i++) {</pre>
  P cur = poly[i], prv = i ? poly[i-1] : poly.back();
  bool side = ori(s, e, cur) < 0;</pre>
  if (side != (ori(s, e, prv) < 0))
   res.push_back(intersect({s, e}, {cur, prv}));
  if (side)
                                                                 sign1) {
   res.push_back(cur);
                                                             PTF v = (b.o - a.o) / d;
 return res:
}
                                                             if (c * c > 1) return {};
6.11 Point In Simple Polygon [037c52]
bool PIP(const vector<P> &p, P z, bool strict = true) {
 int cnt = 0, n = (int)p.size();
for (int i = 0; i < n; i++) {</pre>
  P A = p[i], B = p[(i + 1) \% n];
  if (isInter(Seg(A, B), z)) return !strict;
  auto zy = IM(z), Ay = IM(A), By = IM(B);
                                                              PTF p1 = a.o + n * a.r;
  cnt ^= ((zy<Ay) - (zy<By)) * ori(z, A, B) > 0;
 }
 return cnt;
}
                                                             return ret;
                                                            }
6.12 Point In Hull (Fast) [060ba1]
bool PIH(const vector<P> &h, P z, bool strict = true) {
 int n = (int)h.size(), a = 1, b = n - 1, r = !strict;
 if (n < 3) return r && isInter(Seg(h[0], h[n-1]), z);</pre>
 if (ori(h[0],h[a],h[b]) > 0) swap(a, b);
 if (ori(h[0],h[a],z) >= r || ori(h[0],h[b],z) <= -r)</pre>
                                                             llf dis = abs(o - ft);
  return false;
                                                             if (dis > r) return {};
 while (abs(a - b) > 1) {
  int c = (a + b) / 2;
  (ori(h[0], h[c], z) > 0 ? b : a) = c;
 return ori(h[a], h[b], z) < r;
6.13 Tangent of Points To Hull [6d7cd7]
pair<int, int> get_tangent(const vector<P> &v, P p) {
 const auto gao = [&, N = int(v.size())](int s) {
  const auto lt = [&](int x, int y) {
                                                             llf S, h, theta;
   return ori(p, v[x % N], v[y % N]) == s; };
  int l = 0, r = N; bool up = lt(0, 1);
  while (r - l > 1) {
   int m = (l + r) / 2;
                                                             if (a > r) {
   if (lt(m, 0) ? up : !lt(m, m+1)) r = m;
   else l = m;
  }
                                                              if (h < r && B < PI / 2)
  return (lt(l, r) ? r : l) % N;
                                                             } else if (b > r) {
 }; // test @ codeforces.com/gym/101201/problem/E
 return {gao(-1), gao(1)}; // (a,b):ori(p,v[a],v[b])<0
} // plz ensure that point strictly out of hull
                                                             } else
6.14 Circle Class & Intersection [5111af]
llf FMOD(llf x) {
                                                             return S;
 if (x < -PI) x += PI * 2;
 if (x > PI) x -= PI * 2;
 return x;
                                                             llf S = 0;
struct Cir { PTF o; llf r; };
// be carefule when tangent
vector<llf> intersectAngle(Cir a, Cir b) {
 PTF dir = b.o - a.o; llf d2 = norm(dir);
                                                             return abs(S);
 if (norm(a.r - b.r) >= d2) { // <math>norm(x) := |x|^2
  if (a.r < b.r) return \{-PI, PI\}; // a in b
  else return {}; // b in a
 } else if (norm(a.r + b.r) <= d2) return {};</pre>
 llf dis = abs(dir), theta = arg(dir);
 llf phi = acos((a.r * a.r + d2 - b.r * b.r) /
```

(2 \* a.r \* dis)); // is acos\_safe needed ?

return { L, R };

llf L = FMOD(theta - phi), R = FMOD(theta + phi);

```
if (d > b.r+a.r || d < abs(b.r-a.r)) return {};</pre>
 llf dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
 PTF dir = (a.o - b.o) / d;
 PTF v = rot90(dir) * sqrt(max(0.0L, b.r*b.r-d1*d1));
6.15 Circle Common Tangent [5ff02c]
// be careful of tangent / exact same circle
// sign1 = 1 for outer tang, -1 for inter tang
vector<Line> common_tan(const Cir &a, const Cir &b, int
 if (norm(a.o - b.o) < eps) return {};</pre>
 llf d = abs(a.o - b.o), c = (a.r - sign1 * b.r) / d;
 if (abs(c * c - 1) < eps) {
 PTF p = a.o + c * v * a.r
  return {Line(p, p + rot90(b.o - a.o))};
 vector<Line> ret; llf h = sqrt(max(0.0L, 1-c*c));
 for (int sign2 : {1, -1}) {
 PTF n = c * v + sign2 * h * rot90(v);
  PTF p2 = b.o + n * (b.r * sign1);
  ret.emplace_back(p1, p2);
6.16 Line-Circle Intersection [12b42a]
vector<PTF> LineCircleInter(PTF p1, PTF p2, PTF o, llf
 PTF ft = p1 + project(o-p1, p2-p1), vec = p2-p1;
 if (abs(dis - r) < eps) return {ft};</pre>
 vec = vec * sqrt(r * r - dis * dis) / abs(vec);
 return {ft + vec, ft - vec}; // sqrt_safe?
6.17 Poly-Circle Intersection [7f140a]
// Divides into multiple triangle, and sum up
// from 8BQube, test by HDU2892 & AOJ CGL_7_H
llf _area(PTF pa, PTF pb, llf r) {
 if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
 if (abs(pb) < eps) return 0;</pre>
 llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
 llf cB = dot(pb, pb-pa) / a / c, B = acos_safe(cB);
 llf cC = dot(pa, pb) / a / b, C = acos_safe(cC);
  S = (C / 2) * r * r; h = a * b * sin(C) / c;
   S = (acos\_safe(h/r)*r*r - h*sqrt\_safe(r*r-h*h));
  theta = PI - B - asin_safe(sin(B) / r * a);
  S = 0.5 * a*r*sin(theta) + (C-theta)/2 * r * r;
  S = 0.5 * sin(C) * a * b;
llf area_poly_circle(const vector<PTF> &v, PTF 0, llf r
 for (size_t i = 0, N = v.size(); i < N; ++i)</pre>
  S += _area(v[i] - 0, v[(i + 1) % N] - 0, r) *
     ori(0, v[i], v[(i + 1) % N]);
6.18 Minimum Covering Circle [faa85a]
Cir getCircum(P a, P b, P c){ // P = complex<llf>
 P z 1 = a - b, z 2 = a - c; llf D = cross(z 1, z 2) * 2;
llf c1 = dot(a + b, z1), c2 = dot(a + c, z2);
 P \circ = rot90(c2 * z1 - c1 * z2) / D;
 return { o, abs(o - a) };
```

Cir minCircleCover(vector<P> pts) {

```
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 assert (!pts.empty());
 ranges::shuffle(pts, mt19937(114514));
 Cir c = \{ 0, 0 \};
 for(size_t i = 0; i < pts.size(); i++) {
   if (abs(pts[i] - c.o) <= c.r) continue;</pre>
  c = { pts[i], 0 };
  for (size_t j = 0; j < i; j++) {</pre>
   if (abs(pts[j] - c.o) <= c.r) continue;</pre>
   c.o = (pts[i] + pts[j]) / llf(2);
   c.r = abs(pts[i] - c.o);
   for (size_t k = 0; k < j; k++) {</pre>
    if (abs(pts[k] - c.o) <= c.r) continue;</pre>
    c = getCircum(pts[i], pts[j], pts[k]);
  }
 return c;
} // test @ TIOJ 1093 & luogu P1742
6.19 Circle Union [1a5265]
#define eb emplace_back
struct Teve { // test@SPOJ N=1000, 0.3~0.5s
PTF p; llf a; int add; // point, ang, add
Teve(PTF x, llf y, int z) : p(x), a(y), add(z) {}
 bool operator<(Teve &b) const { return a < b.a; }</pre>
// strict: x = 0, otherwise x = -1
bool disjunct(Cir &a, Cir &b, int x)
{ return sgn(abs(a.o - b.o) - a.r - b.r) > x; }
bool contain(Cir &a, Cir &b, int x)
{ return sgn(a.r - b.r - abs(a.o - b.o)) > x; }
vector<llf> CircleUnion(vector<Cir> &c) {
 // area[i] : area covered by at least i circles
 int N = (int)c.size(); vector<llf> area(N + 1);
 vector<vector<int>> overlap(N, vector<int>(N));
 auto g = overlap; // use simple 2darray to speedup
 for (int i = 0; i < N; ++i)</pre>
  for (int j = 0; j < N; ++j) {
   /* c[j] is non-strictly in c[i]. */
   overlap[i][j] = i != j &&
    (sgn(c[i].r - c[j].r) > 0 | |
     (sgn(c[i].r - c[j].r) == 0 && i < j)) &&
    contain(c[i], c[j], -1);
 for (int i = 0; i < N; ++i)
  for (int j = 0; j < N; ++j)
g[i][j] = i != j && !(overlap[i][j] ||</pre>
     overlap[j][i] || disjunct(c[i], c[j], -1));
 for (int i = 0; i < N; ++i) {</pre>
  vector<Teve> eve; int cnt = 1;
  for (int j = 0; j < N; ++j) cnt += overlap[j][i];</pre>
  // if (cnt > 1) continue; (if only need area[1])
  for (int j = 0; j < N; ++j) if (g[i][j]) {</pre>
   auto IP = intersectPoint(c[i], c[j]);
   PTF aa = IP[1], bb = IP[0];
llf A = arg(aa - c[i].o), B = arg(bb - c[i].o);
   eve.eb(bb, B, 1); eve.eb(aa, A, -1);
   if (B > A) ++cnt;
  if (eve.empty()) area[cnt] += PI*c[i].r*c[i].r;
```

# 6.20 Polygon Union [2bff43]

sort(eve.begin(), eve.end());

cnt += eve[j].add;

eve.eb(eve[0]); eve.back().a += PI  $\star$  2; for (size\_t j = 0; j + 1 < eve.size(); j++) {</pre>

llf t = eve[j + 1].a - eve[j].a;

area[cnt] += cross(eve[j].p, eve[j+1].p) \*.5;

area[cnt] += (t-sin(t)) \* c[i].r \* c[i].r \*.5;

else {

} }

return area;

```
llf rat(P a, P b) { return sgn(RE(b)) ? llf(RE(a))/RE(b
    ) : llf(IM(a))/IM(b); }
llf polyUnion(vector<vector<P>>& poly) {
llf ret = 0; // area of poly[i] must be non-negative
rep(i,0,sz(poly)) rep(v,0,sz(poly[i])) {
 P A = poly[i][v], B = poly[i][(v + 1) % sz(poly[i])];
 vector<pair<llf, int>> segs{{0, 0}, {1, 0}};
```

```
rep(j,0,sz(poly)) if (i != j) {
  rep(u,0,sz(poly[j])) {
    P C = poly[j][u], D = poly[j][(u + 1) % sz(poly[j])
    ];
    if (int sc = ori(A, B, C), sd = ori(A, B, D); sc !=
     sd) {
     llf sa = cross(D-C, A-C), sb = cross(D-C, B-C);
    if (min(sc, sd) < 0)
     segs.emplace_back(sa / (sa - sb), sgn(sc - sd));
    } else if (!sc && !sd && j<i && sgn(dot(B-A,D-C))</pre>
    >0){
    segs.emplace_back(rat(C - A, B - A), 1);
    segs.emplace_back(rat(D - A, B - A), -1);
  }
  sort(segs.begin(), segs.end());
  for (auto &s : segs) s.first = clamp<llf>(s.first, 0,
    1);
  llf sum = 0;
  int cnt = segs[0].second;
  rep(j,1,sz(segs)) {
   if (!cnt) sum += segs[j].first - segs[j - 1].first;
  cnt += segs[j].second;
 ret += cross(A,B) * sum;
}
 return ret / 2;
6.21 3D Point [46b73b]
struct P3 {
```

```
lld x, y, z;
 P3 operator^(const P3 &b) const {
  return {y*b.z-b.y*z, z*b.x-b.z*x, x*b.y-b.x*y};
 //Azimuthal angle (longitude) to x-axis. \in [-pi, pi]
 llf phi() const { return atan2(y, x); }
 //Zenith angle (latitude) to the z-axis. \in [0, pi]
 llf theta() const { return atan2(sqrt(x*x+y*y),z); }
P3 ver(P3 a, P3 b, P3 c) { return (b - a) ^ (c - a); }
lld volume(P3 a, P3 b, P3 c, P3 d) {
 return dot(ver(a, b, c), d - a);
P3 rotate_around(P3 p, llf angle, P3 axis) {
 llf s = sin(angle), c = cos(angle);
 P3 u = normalize(axis);
 return u*dot(u, p)*(1-c) + p * c + cross(u, p)*s;
```

#### 6.22 3D Convex Hull [01652a]

```
struct Face {
 int a, b, c;
 Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
auto preprocess(const vector<P3> &pt) {
 auto G = pt.begin();
 auto a = find_if(all(pt), [&](P3 z) {
 return z != *G; }) - G;
 auto b = find_if(all(pt), [&](P3 z) {
  return ver(*G, pt[a], z) != P3(0, 0, 0); }) - G;
 auto c = find_if(all(pt), [&](P3 z) {
  return volume(*G, pt[a], pt[b], z) != 0; }) - G;
 vector<size_t> id;
 for (size_t i = 0; i < pt.size(); i++)</pre>
  if (i != a && i != b && i != c) id.push_back(i);
 return tuple{a, b, c, id};
// return the faces with pt indexes
// all points coplanar case will WA
vector<Face> convex_hull_3D(const vector<P3> &pt) {
 const int n = int(pt.size());
 if (n <= 3) return {}; // be careful about edge case</pre>
 vector<Face> now;
 vector<vector<int>> z(n, vector<int>(n));
 auto [a, b, c, ord] = preprocess(pt);
 now.emplace_back(a, b, c); now.emplace_back(c, b, a);
 for (auto i : ord) {
  vector<Face> next;
  for (const auto &f : now) {
   lld v = volume(pt[f.a], pt[f.b], pt[f.c], pt[i]);
```

```
if (v <= 0) next.push_back(f);</pre>
  z[f.a][f.b] = z[f.b][f.c] = z[f.c][f.a] = sgn(v);
  const auto F = [\&](int x, int y) \{
   if (z[x][y] > 0 && z[y][x] <= 0)
   next.emplace_back(x, y, i);
 for (const auto &f : now)
  F(f.a, f.b), F(f.b, f.c), F(f.c, f.a);
 now = next;
return now;
}
// n^2 delaunay: facets with negative z normal of
// convexhull of (x, y, x^2 + y^2), use a pseudo-point
// (0, 0, inf) to avoid degenerate case
// test @ SPOJ CH3D
// llf area = 0, vol = 0; // surface area / volume
// for (auto [a, b, c]: faces)
   area += abs(ver(p[a], p[b], p[c]))/2.0,
// vol += volume(P3(0, 0, 0), p[a], p[b], p[c])/6.0;
```

# **6.23 3D Projection** [68f350]

```
using P3F = valarray<llf>;
P3F toP3F(P3 p) { return {p.x, p.y, p.z}; }
llf dot(P3F a, P3F b) {
  return a[0]*b[0]+a[1]*b[1]+a[2]*b[2];
}
P3F housev(P3 A, P3 B, int s) {
  const llf a = abs(A), b = abs(B);
  return toP3F(A) / a + s * toP3F(B) / b;
}
P project(P3 p, P3 q) {
  P3 o(0, 0, 1);
  P3F u = housev(q, o, q.z > 0 ? 1 : -1);
  auto pf = toP3F(p);
  auto np = pf - 2 * u * dot(u, pf) / dot(u, u);
  return P(np[0], np[1]);
} // project p onto the plane q^Tx = 0
```

#### 6.24 3D Skew Line Nearest Point

```
 \begin{array}{l} \boldsymbol{\cdot} \  \, L_1: \boldsymbol{v}_1 = \boldsymbol{p}_1 + t_1 \boldsymbol{d}_1, L_2: \boldsymbol{v}_2 = \boldsymbol{p}_2 + t_2 \boldsymbol{d}_2 \\ \boldsymbol{\cdot} \  \, \boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2 \\ \boldsymbol{\cdot} \  \, \boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}, \boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n} \\ \boldsymbol{\cdot} \  \, \boldsymbol{c}_1 = \boldsymbol{p}_1 + \frac{(\boldsymbol{p}_2 - \boldsymbol{p}_1) \cdot \boldsymbol{n}_2}{\boldsymbol{d}_1 \cdot \boldsymbol{n}_2} \boldsymbol{d}_1, \boldsymbol{c}_2 = \boldsymbol{p}_2 + \frac{(\boldsymbol{p}_1 - \boldsymbol{p}_2) \cdot \boldsymbol{n}_1}{\boldsymbol{d}_2 \cdot \boldsymbol{n}_1} \boldsymbol{d}_2 \end{array}
```

#### 6.25 Delaunay [3a4ff1]

```
/* please ensure input points are unique */
/* A triangulation such that no points will strictly
inside circumcircle of any triangle. C should be big
enough s.t. the initial triangle contains all points */
#define L(i) ((i)==0 ? 2 : (i)-1)
#define R(i) ((i)==2 ? 0 : (i)+1)
#define F3 for (int i = 0; i < 3; i++)
bool is_inf(P z) { return RE(z) <= -C || RE(z) >= C; }
bool in_cc(const array<P,3> &p, P q) {
 i128 inf_det = 0, det = 0, inf_N, N;
 F3 {
  if (is_inf(p[i]) && is_inf(q)) continue;
  else if (is_inf(p[i])) inf_N = 1, N = -norm(q);
 else if (is_inf(q)) inf_N = -1, N = norm(p[i]);
else inf_N = 0, N = norm(p[i]) - norm(q);
  lld D = cross(p[R(i)] - q, p[L(i)] - q);
  inf_det += inf_N * D; det += N * D;
 return inf_det != 0 ? inf_det > 0 : det > 0;
P v[maxn];
struct Tri;
struct E {
 Tri *t; int side;
E(Tri *t_=0, int side_=0) : t(t_), side(side_) {}
struct Tri {
 array<int,3> p; array<Tri*,3> ch; array<E,3> e;
 Tri(int a=0, int b=0, int c=0) : p{a, b, c}, ch{} {}
 bool has_chd() const { return ch[0] != nullptr; }
 bool contains(int q) const {
 F3 if (ori(v[p[i]], v[p[R(i)]], v[q]) < 0)
   return false;
  return true;
```

```
bool check(int q) const {
  return in_cc({v[p[0]], v[p[1]], v[p[2]]}, v[q]); }
} pool[maxn * 10], *it, *root;
void link(const E &a, const E &b) {
if (a.t) a.t->e[a.side] = b;
 if (b.t) b.t->e[b.side] = a;
void flip(Tri *A, int a) {
 auto [B, b] = A->e[a]; /* flip edge between A,B */
 if (!B || !A->check(B->p[b])) return;
 Tri *X = new (it++) Tri(A->p[R(a)], B->p[b], A->p[a]);
 Tri \star Y = new (it++) Tri(B->p[R(b)], A->p[a], B->p[b]);
 link(E(X, 0), E(Y, 0));
 link(E(X, 1), A\rightarrow e[L(a)]); link(E(X, 2), B\rightarrow e[R(b)]);
 link(E(Y, 1), B->e[L(b)]); link(E(Y, 2), A->e[R(a)]);
 A->ch = B->ch = {X, Y, nullptr};
flip(X, 1); flip(X, 2); flip(Y, 1); flip(Y, 2);
void add_point(int p) {
 Tri *r = root;
 while (r->has_chd()) for (Tri *c: r->ch)
  if (c && c->contains(p)) { r = c; break; }
 array<Tri*, 3> t; /* split into 3 triangles */
 F3 t[i] = new (it++) Tri(r->p[i], r->p[R(i)], p);
 F3 link(E(t[i], 0), E(t[R(i)], 1));
F3 link(E(t[i], 2), r->e[L(i)]);
 r->ch = t;
 F3 flip(t[i], 2);
auto build(const vector<P> &p) {
 it = pool; int n = (int)p.size();
 vector<int> ord(n); iota(all(ord), 0);
 shuffle(all(ord), mt19937(114514));
 root = new (it++) Tri(n, n + 1, n + 2);
 copy_n(p.data(), n, v); v[n++] = P(-C, -C);
 v[n++] = P(C * 2, -C); v[n++] = P(-C, C * 2);
 for (int i : ord) add_point(i);
 vector<array<int, 3>> res;
 for (Tri *now = pool; now != it; now++)
  if (!now->has_chd()) res.push_back(now->p);
 return res;
```

#### 6.26 Build Voronoi [94f000]

```
void build_voronoi_cells(auto &&p, auto &&res) {
  vector<vector<int>> adj(p.size());
  for (auto f: res) F3 {
    int a = f[i], b = f[R(i)];
    if (a >= p.size() || b >= p.size()) continue;
    adj[a].emplace_back(b);
}
// use `adj` and `p` and HPI to build cells
  for (size_t i = 0; i < p.size(); i++) {
    vector<Line> ls = frame; // the frame
    for (int j : adj[i]) {
        P m = p[i] + p[j], d = rot90(p[j] - p[i]);
        assert (norm(d) != 0);
        ls.emplace_back(m, m + d); // doubled coordinate
    } // HPI(ls)
}
```

#### 6.27 kd Tree (Nearest Point) [dbade8]

```
struct KDTree {
 struct Node {
  int x, y, x1, y1, x2, y2, id, f; Node *L, *R;
 } tree[maxn], *root;
 lld dis2(int x1, int y1, int x2, int y2) {
 lld dx = x1 - x2, dy = y1 - y2;
  return dx * dx + dy * dy;
 static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
 static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 void init(vector<pair<int,int>> &ip) {
  const int n = ip.size();
  for (int i = 0; i < n; i++) {</pre>
   tree[i].id = i:
   tree[i].x = ip[i].first;
   tree[i].y = ip[i].second;
  }
  root = build(0, n-1, 0);
```

```
Node* build(int L, int R, int d) {
  if (L>R) return nullptr; int M = (L+R)/2;
  nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
  Node &o = tree[M]; o.f = d % 2;
o.x1 = o.x2 = o.x; o.y1 = o.y1 = o.y;
  o.L = build(L, M-1, d+1); o.R = build(M+1, R, d+1);
  for (Node *s: {o.L, o.R}) if (s) {
   o.x1 = min(o.x1, s->x1); o.x2 = max(o.x2, s->x2);
   o.y1 = min(o.y1, s->y1); o.y2 = max(o.y2, s->y2);
  }
  return tree+M;
 bool touch(int x, int y, lld d2, Node *r){
  lld d = sqrt(d2)+1;
  return x >= r->x1 - d && x <= r->x2 + d &&
         y >= r->y1 - d \&\& y <= r->y2 + d;
 using P = pair<lld, int>;
 void dfs(int x, int y, P &mn, Node *r) {
  if (!r || !touch(x, y, mn.first, r)) return;
  mn = min(mn, P(dis2(r->x, r->y, x, y), r->id));
  if (r->f == 1 ? y < r->y : x < r->x)
   dfs(x, y, mn, r\rightarrow L), dfs(x, y, mn, r\rightarrow R);
  else
   dfs(x, y, mn, r\rightarrow R), dfs(x, y, mn, r\rightarrow L);
 int query(int x, int y) {
  P mn(INF, -1); dfs(x, y, mn, root);
  return mn.second;
} tree;
```

# 6.28 kd Closest Pair (3D ver.) [84d9eb]

```
llf solve(vector<P> v) {
shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,</pre>
  unordered_map<lld, int>>> m;
llf d = dis(v[0], v[1]);
auto Idx = [\&d] (llf x) \rightarrow lld {
  return round(x * 2 / d) + 0.1; };
 auto rebuild_m = [&m, &v, &Idx](int k) {
  m.clear();
  for (int i = 0; i < k; ++i)
   m[Idx(v[i].x)][Idx(v[i].y)]
    \lceil Idx(v[i].z) \rceil = i;
}; rebuild_m(2);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
     kz = Idx(v[i].z); bool found = false;
  for (int dx = -2; dx <= 2; ++dx) {
   const lld nx = dx + kx;
   if (m.find(nx) == m.end()) continue;
   auto& mm = m[nx];
   for (int dy = -2; dy <= 2; ++dy) {
    const lld ny = dy + ky;
if (mm.find(ny) == mm.end()) continue;
    auto& mmm = mm[ny];
    for (int dz = -2; dz <= 2; ++dz) {
  const lld nz = dz + kz;</pre>
     if (mmm.find(nz) == mmm.end()) continue;
     const int p = mmm[nz];
     if (dis(v[p], v[i]) < d) {
      d = dis(v[p], v[i]);
      found = true;
   }
  if (found) rebuild_m(i + 1);
  else m[kx][ky][kz] = i;
return d;
```

# 6.29 Simulated Annealing [4e0fe5]

```
llf anneal() {
  mt19937 rnd_engine(seed);
  uniform_real_distribution<llf> rnd(0, 1);
  const llf dT = 0.001;
  // Argument p
  llf S_cur = calc(p), S_best = S_cur;
  for (llf T = 2000; T > EPS; T -= dT) {
```

```
// Modify p to p_prime
  const llf S_prime = calc(p_prime);
  const llf delta_c = S_prime - S_cur;
  llf prob = min((llf)1, exp(-delta_c / T));
if (rnd(rnd_engine) <= prob)</pre>
   S_cur = S_prime, p = p_prime;
  if (S_prime < S_best) // find min</pre>
   S_best = S_prime, p_best = p_prime;
 return S_best;
6.30 Triangle Centers [adb146]
0 = ... // see min circle cover
G = (A + B + C) / 3;
H = G * 3 - 0 * 2; // orthogonal center
llf a = abs(B - C), b = abs(A - C), c = abs(A - B);
I = (a * A + b * B + c * C) / (a + b + c);
// FermatPoint: minimizes sum of distance
// if max. angle >= 120 deg then vertex
// otherwise, make eq. triangle AB'C, CA'B, BC'A
// line AA', BB', CC' intersects at P
     Stringology
     Hash [ce7fad]
template <int P = 127, int Q = 1051762951>
class Hash {
 vector<int> h, p;
public:
 Hash(const auto &s) : h(s.size()+1), p(s.size()+1) {
   for (size_t i = 0; i < s.size(); ++i)</pre>
   h[i + 1] = add(mul(h[i], P), s[i]);
  generate(all(p), [x = 1, y = 1, this]() mutable {
   return y = x, x = mul(x, P), y; y; y;
 int query(int l, int r) const { // 1-base (l, r]
  return sub(h[r], mul(h[l], p[r - l]));
 }
};
7.2 Suffix Array [1f4d4f]
namespace sfx {
bool _t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], _c[maxn * 2];
int x[maxn], _p[maxn], _q[maxn * 2];
// sa[i]: sa[i]-th suffix is the
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *a, int *c, int n, int z) {
 memset(a, 0, sizeof(int) * n);
 memcpy(x, c, sizeof(int) * z);
void induce(int *a, int *c, int *s,
  bool *t, int n, int z) {
 memcpy(x + 1, c, sizeof(int) * (z - 1));
 for (int i = 0; i < n; ++i)
  if (a[i] && !t[a[i] - 1])
   a[x[s[a[i] - 1]]++] = a[i] - 1;
 memcpy(x, c, sizeof(int) * z);
for (int i = n - 1; i >= 0; --
  if (a[i] && t[a[i] - 1])
   a[--x[s[a[i] - 1]]] = a[i] - 1;
void sais(int *s, int *a, int *p, int *q,
bool *t, int *c, int n, int z) {
bool uniq = t[n - 1] = true;
int nn=0, nz=-1, *nsa = a+n, *ns=s+n, last=-1;
 memset(c, 0, sizeof(int) * z);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];</pre>
 if (uniq) {
  for (int i = 0; i < n; ++i) a[--c[s[i]]] = i;</pre>
 for (int i = n - 2; i >= 0; --i)
t[i] = (s[i]==s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);</pre>
 pre(a, c, n, z);
 for (int i = 1; i <= n - 1; ++i)</pre>
```

if (t[i] && !t[i - 1])

a[--x[s[i]]] = p[q[i] = nn++] = i;

```
induce(a, c, s, t, n, z);

for (int i = 0; i < n; ++i)
  if (a[i] && t[a[i]] && !t[a[i] - 1]) {
   bool neq = last < 0 || memcmp(s + a[i], s + last,</pre>
     (p[q[a[i]] + 1] - a[i]) * sizeof(int));
   ns[q[last = a[i]]] = nz += neq;
 sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nz+1);
 pre(a, c, n, z);

for (int i = nn - 1; i >= 0; --i)
  a[--x[s[p[nsa[i]]]]] = p[nsa[i]];
 induce(a, c, s, t, n, z);
void build(const string &s) {
 const int n = int(s.size());
 for (int i = 0; i < n; ++i) _s[i] = s[i];</pre>
 _s[n] = 0; // s shouldn't contain 0
 sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
for(int i = 0; i < n; ++i) rev[sa[i] = sa[i+1]] = i;</pre>
 int ind = hi[0] = 0;
 for (int i = 0; i < n; ++i) {</pre>
  if (!rev[i]) { ind = 0; continue; }
  while (i + ind < n &&</pre>
    s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  hi[rev[i]] = ind ? ind-- : 0;
}}
```

```
Ex SAM [58374b]
int len[maxn * 2], link[maxn * 2]; // maxlen, suflink
int next[maxn * 2][maxc], tot; // [0, tot), root = 0
int ord[maxn * 2]; // topo. order (sort by length)
int cnt[maxn * 2]; // occurence
int newnode() {
 fill_n(next[tot], maxc, 0);
 return len[tot] = cnt[tot] = link[tot] = 0, tot++;
void init() { tot = 0, newnode(), link[0] = -1; }
int insertSAM(int last, int c) {
 int cur = next[last][c];
 len[cur] = len[last] + 1;
 int p = link[last];
 while (p != -1 && !next[p][c])
  next[p][c] = cur, p = link[p];
 if (p == -1) return link[cur] = 0, cur;
 int q = next[p][c];
 if (len[p] + 1 == len[q]) return link[cur] = q, cur;
 int clone = newnode();
 for (int i = 0; i < maxc; ++i)</pre>
  next[clone][i] = len[next[q][i]] ? next[q][i] : 0;
 len[clone] = len[p] + 1;
 while (p != -1 && next[p][c] == q)
  next[p][c] = clone, p = link[p];
 link[link[cur] = clone] = link[q];
 link[q] = clone;
 return cur;
}
void insert(const string &s) {
 int cur = 0;
 for (char ch : s) {
  int &nxt = next[cur][int(ch - 'a')];
  if (!nxt) nxt = newnode();
  cnt[cur = nxt] += 1;
 }
}
void build() {
 queue<int> q; q.push(0);
 while (!q.empty()) {
  int cur = q.front(); q.pop();
  for (int i = 0; i < maxc; ++i)</pre>
   if (next[cur][i]) q.push(insertSAM(cur, i));
 vector<int> lc(tot);
 for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
 partial_sum(all((lc), lc.begin());
for (int i = 1; i < tot; ++i) ord[--lc[len[i]]] = i;</pre>
void solve() {
 for (int i = tot - 2; i >= 0; --i)
  cnt[link[ord[i]]] += cnt[ord[i]];
```

```
};
7.4 Z value [6a7fd0]
```

```
vector<int> Zalgo(const string &s) {
 vector<int> z(s.size(), s.size());
 for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
  int j = clamp(r - i, 0, z[i - l]);
  for (; i + j < z[0] and s[i + j] == s[j]; ++j);
 if (i + (z[i] = j) > r) r = i + z[l = i];
return z;
```

# 7.5 Manacher [c938a9]

```
vector<int> manacher(const string &S) {
 const int n = (int)S.size(), m = n * 2 + 1;
 vector<int> z(m);
 string t = "."; for (char c: S) t += c, t += '.';
 for (int i = 1, l = 0, r = 0; i < m; ++i) {
  z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
  while (i - z[i] >= 0 \&\& i + z[i] < m) {
   if (t[i - z[i]] == t[i + z[i]]) ++z[i];
   else break;
  if (i + z[i] > r) r = i + z[i], l = i;
 return z; // the palindrome lengths are z[i] - 1
/* for (int i = 1; i + 1 < m; ++i) {
  int l = (i - z[i] + 2) / 2, r = (i + z[i]) / 2;
  if (l != r) // [l, r) is maximal palindrome
```

# 7.6 Lyndon Factorization [d22cc9]

```
// partition s = w[0] + w[1] + ... + w[k-1],
// w[0] >= w[1] >= ... >= w[k-1]
// each w[i] strictly smaller than all its suffix
void duval(const auto &s, auto &&report) {
 for (int n = (int)s.size(), i = 0, j, k; i < n; ) {</pre>
  for (j = i + 1, k = i; j < n \&\& s[k] <= s[j]; j++)
   k = (s[k] < s[j] ? i : k + 1);
  // if (i < n / 2 && j >= n / 2) {
  // for min cyclic shift, call duval(s + s)
  // then here s.substr(i, n / 2) is min cyclic shift
  // }
  for (; i <= k; i += j - k)
   report(i, j - k); // s.substr(l, len)
} // tested @ luogu 6114, 1368 & UVA 719
```

# 7.7 Main Lorentz [615b8f]

```
vector<pair<int, int>> rep[kN]; // 0-base [l, r]
void main_lorentz(const string &s, int sft = 0) {
 const int n = s.size();
 if (n == 1) return;
 const int nu = n / 2, nv = n - nu;
 const string u = s.substr(0, nu), v = s.substr(nu),
    ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend());
main_lorentz(u, sft), main_lorentz(v, sft + nu);
const auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u),
    z3 = Zalgo(ru + '#' + rv), z4 = Zalgo(v);
 auto get_z = [](const vector<int> &z, int i) {
  return (0 <= i and i < (int)z.size()) ? z[i] : 0; };</pre>
 auto add_rep = [&](bool left, int c, int l, int k1,
    int k2) {
  const int L = max(1, l - k2), R = min(l - left, k1);
  if (L > R) return;
  if (left) rep[l].emplace_back(sft + c - R, sft + c -
    L);
  else rep[l].emplace_back(sft + c - R - l + 1, sft + c
      - L - l + 1);
 for (int cntr = 0; cntr < n; cntr++) {</pre>
  int l, k1, k2;
  if (cntr < nu) {</pre>
   l = nu - cntr;
   k1 = get_z(z1, nu - cntr);
k2 = get_z(z2, nv + 1 + cntr);
  } else {
   l = cntr - nu + 1;
   k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
   k2 = get_z(z4, (cntr - nu) + 1);
```

```
if (k1 + k2 >= l)
   add_rep(cntr < nu, cntr, l, k1, k2);</pre>
}
```

#### 7.8 BWT [5a9b3a]

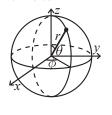
```
vector<int> v[SIGMA];
void BWT(char *ori, char *res) {
  // make ori -> ori + ori
// then build suffix array
void iBWT(char *ori, char *res) {
for (int i = 0; i < SIGMA; i++) v[i].clear();</pre>
const int len = strlen(ori);
for (int i = 0; i < len; i++)</pre>
 v[ori[i] - 'a'].push_back(i);
vector<int> a;
 for (int i = 0, ptr = 0; i < SIGMA; i++)</pre>
 for (int j : v[i]) {
   a.push_back(j);
ori[ptr++] = 'a' + i;
for (int i = 0, ptr = 0; i < len; i++) {</pre>
  res[i] = ori[a[ptr]];
  ptr = a[ptr];
res[len] = 0;
```

#### 7.9 Palindromic Tree [0673ee]

```
struct PalindromicTree {
 struct node {
  int nxt[26], f, len; // num = depth of fail link
  int cnt, num;  // = #pal_suffix of this node
node(int l = 0) : nxt{}, f(0), len(l), cnt(0), num(0)
 vector<node> st; vector<char> s; int last, n;
 void init() {
  st.clear(); s.clear();
  last = 1; n = 0;
  st.push_back(0); st.push_back(-1);
  st[0].f = 1; s.push_back(-1);
 int getFail(int x) {
  while (s[n - st[x].len - 1] != s[n]) x = st[x].f;
  return x;
 void add(int c) {
  s.push_back(c -= 'a'); ++n;
  int cur = getFail(last);
  if (!st[cur].nxt[c]) {
   int now = st.size();
   st.push_back(st[cur].len + 2);
   st[now].f = st[getFail(st[cur].f)].nxt[c];
   st[cur].nxt[c] = now;
   st[now].num = st[st[now].f].num + 1;
  last = st[cur].nxt[c]; ++st[last].cnt;
 void dpcnt() { // cnt = #occurence in whole str
  for (int i = st.size() - 1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
 int size() { return st.size() - 2; }
} pt;
/* usage
string s; cin >> s; pt.init();
for (int i = 0; i < SZ(s); i++) {
 int prvsz = pt.size(); pt.add(s[i]);
 if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
  // pal @ [l,r]: s.substr(l, r-l+1)
} */
```

# 8 Misc 8.1 Theorems

# **Spherical Coordinate**



$$\begin{split} r &= \sqrt{x^2 + y^2 + z^2} \\ \theta &= \mathsf{acos}(z/\sqrt{x^2 + y^2 + z^2}) \\ \phi &= \mathsf{atan2}(y,x) \end{split}$$

#### Sherman-Morrison formula

$$(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1+v^{\mathsf{T}}A^{-1}u}$$

#### Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} =$ d(i),  $L_{ij}=-c$  where c is the number of edge (i,j) in  $\tilde{G}$ 

- The number of undirected spanning in G is  $\det(\tilde{L}_{11})$ .
- The number of directed spanning tree rooted at r in G is  $\det(\tilde{L}_{rr})$ .

#### Tutte's Matrix

 $x=r\sin\theta\cos\phi$ 

 $y = r \sin \theta \sin \phi$ 

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

# Cayley's Formula

- Given a degree sequence  $d_1, d_2, \ldots, d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\dots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .

#### Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all  $1 \le k \le n$ .

#### Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other great-

#### Euler's planar graph formula

$$V - E + F = C + 1$$
.  $E \le 3V - 6$  (when  $V \ge 3$ )

For simple polygon, when points are all integer, we have  $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$ 

#### Matroid Intersection

Given matroids  $M_1=(G,I_1),M_2=(G,I_2)$ , find maximum  $S\in I_1\cap I_2$ . For each iteration, build the directed graph and find a shortest path from s to t.

```
• s \to x : S \sqcup \{x\} \in I_1
• x \to t : S \sqcup \{x\} \in I_2
```

 $\begin{array}{lll} \cdot & y \to x: S \setminus \{x\} \cup \{x\} \in I_2 & \text{is in the unique circuit of } S \cup \{x\}\} \\ \cdot & x \to y: S \setminus \{y\} \cup \{x\} \in I_2 & \text{is in the unique circuit of } S \cup \{x\}\} \\ \text{Alternate the path, and } & |S| & \text{will increase by 1.} & \text{Let } R & = \\ \min(\text{rank}(I_1), \text{rank}(I_2)), & N & = & |G|. & \text{In each iteration, } |E| & = & O(RN). \\ \text{For weighted case, assign weight } & -w(x) & \text{and } w(x) & \text{to } x \in S & \text{and } x \notin S, \\ \end{array}$ resp. Use Bellman-Ford to find the weighted shortest path. The maximum

# iteration of Bellman-Ford is 2R+1. **8.2 Weight Matroid Intersection** [d00ee8]

```
struct Matroid {
 Matroid(bitset<N>); // init from an independent set
bool can_add(int); // check if break independence
Matroid remove(int); // removing from the set
auto matroid_intersection(const vector<int> &w) {
 const int n = (int)w.size(); bitset<N> S;
 for (int sz = 1; sz <= n; sz++) {</pre>
  Matroid M1(S), M2(S); vector<vector<pii>>> e(n + 2);
  for (int j = 0; j < n; j++) if (!S[j]) {</pre>
   if (M1.can_add(j)) e[n].eb(j, -w[j]);
if (M2.can_add(j)) e[j].eb(n + 1, 0);
  for (int i = 0; i < n; i++) if (S[i]) {</pre>
   Matroid T1 = M1.remove(i), T2 = M2.remove(i);
    for (int j = 0; j < n; j++) if (!S[j]) {</pre>
```

```
if (T1.can_add(j)) e[i].eb(j, -w[j]);
    if (T2.can_add(j)) e[j].eb(i, w[i]);
   }
  } // maybe implicit build graph for more speed
  vector<pii> d(n + 2, {INF, 0}); d[n] = {0, 0};
  vector<int> prv(n + 2, -1);
  // change to SPFA for more speed, if necessary
  for (int upd = 1; upd--; )
   for (int u = 0; u < n + 2; u++)
    for (auto [v, c] : e[u]) {
  pii x(d[u].first + c, d[u].second + 1);
     if (x < d[v]) d[v] = x, prv[v] = u, upd = 1;
    }
  if (d[n + 1].first >= INF) break;
  for (int x = prv[n+1]; x!=n; x = prv[x]) S.flip(x);
  // S is the max-weighted independent set w/ size sz
 return S;
} // from Nacl
```

# 8.3 Stable Marriage

```
1: Initialize m \in M and w \in W to free
2: while \exists free man m who has a woman w to propose to do
         w \leftarrow \text{first woman on } m \text{'s list to whom } m \text{ has not yet proposed}
        if \exists some pair (m', w) then
            if w prefers m to m^\prime then
6:
7:
                \dot{m}' \leftarrow \textit{free}
                 (m, w) \leftarrow \mathsf{engaged}
8:
            end if
9:
10.
             (m, w) \leftarrow \mathsf{engaged}
11:
        end if
12: end while
```

# 8.4 Bitset LCS [330ab1]

```
cin >> n >> m;
for (int i = 1, x; i <= n; ++i)
  cin >> x, p[x].set(i);
for (int i = 1, x; i <= m; i++) {
  cin >> x, (g = f) |= p[x];
  f.shiftLeftByOne(), f.set(0);
  ((f = g - f) ^= g) &= g;
}
cout << f.count() << '\n';</pre>
```

# 8.5 Prefix Substring LCS [7d8faf]

```
void all_lcs(string S, string T) { // 0-base
vector<size_t> h(T.size()); iota(all(h), 1);
for (size_t a = 0; a < S.size(); ++a) {
  for (size_t c = 0, v = 0; c < T.size(); ++c)
    if (S[a] == T[c] || h[c] < v) swap(h[c], v);
    // here, LCS(s[0, a], t[b, c]) =
    // c - b + 1 - sum([h[i] > b] | i <= c)
  }
} // test @ yosupo judge</pre>
```

#### 8.6 Convex 1D/1D DP [6e0124]

```
struct segment {
 int i, l, r;
 segment() {}
 segment(int a, int b, int c): i(a), l(b), r(c) {}
void solve() {
 auto f = [](int l, int r){return dp[l] + w(l+1, r);}
 deque<segment> dq; dq.push_back(segment(0, 1, n));
 for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i);
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
   f(i, dq.back().l)<f(dq.back().i, dq.back().l))</pre>
    dq.pop_back();
  if (dq.size()) {
   int d = 1 << 20, c = dq.back().l;</pre>
   while (d >>= 1) if (c + d <= dq.back().r)</pre>
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
  if (seg.l <= n) dq.push_back(seg);</pre>
}
```

# 8.7 ConvexHull Optimization [b4318e]

```
struct L {
  mutable lld a, b, p;
  bool operator<(const L &r) const {</pre>
  return a < r.a; /* here */ }
 bool operator<(lld x) const { return p < x; }</pre>
lld Div(lld a, lld b) {
  return a / b - ((a ^ b) < 0 && a % b); }</pre>
 struct DynamicHull : multiset<L, less<>>> {
  static const lld kInf = 1e18;
  bool Isect(iterator x, iterator y) {
   if (y == end()) { x->p = kInf; return false; }
   if (x->a == y->a)
    x->p = x->b > y->b ? kInf : -kInf; /* here */
   else x -> p = Div(y -> b - x -> b, x -> a - y -> a);
   return x->p >= y->p;
  void Insert(lld a, lld b) {
   auto z = insert({a, b, 0}), y = z++, x = y;
   while (Isect(y, z)) z = erase(z);
   if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
   while ((y = x) != begin() && (--x)->p >= y->p)
    Isect(x, erase(y));
  lld Query(lld x) { // default chmax
   auto l = *lower_bound(x); // to chmin:
   return l.a * x + l.b; // modify the 2 "<>"
};
```

#### 8.8 Min Plus Convolution [464dcd]

#### 8.9 De-Bruijn [c0a223]

```
vector<int> de_bruijn(int k, int n) {
    // return cyclic string of len k^n s.t. every string
    // of len n using k char appears as a substring.
    vector<int> aux(n + 1), res;
    auto db = [&](auto self, int t, int p) -> void {
        if (t <= n)
            for (int i = aux[t - p]; i < k; ++i, p = t)
                aux[t] = i, self(self, t + 1, p);
        else if (n % p == 0) for (int i = 1; i <= p; ++i)
        res.push_back(aux[i]);
    }; db(db, 1, 1);
    return res;
}</pre>
```

### 8.10 Josephus Problem [f4494f]

```
int f(int n, int m) { // n people kill m for each turn
  int s = 0;
  for (int i = 2; i <= n; i++) s = (s + m) % i;
  return s;
}
int kth(int n, int m, int k){ // died at kth
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

#### 8.11 N Queens Problem [31f83e]

```
void solve(VI &ret, int n) { // no sol when n=2,3
if (n % 6 == 2) {
  for (int i = 2; i <= n; i += 2) ret.push_back(i);
  ret.push_back(3); ret.push_back(1);
  for (int i = 7; i <= n; i += 2) ret.push_back(i);
  ret.push_back(5);
} else if (n % 6 == 3) {</pre>
```

```
for (int i = 4; i <= n; i += 2) ret.push_back(i);</pre>
  ret.push_back(2);
  for (int i = 5; i <= n; i += 2) ret.push_back(i);</pre>
  ret.push_back(1); ret.push_back(3);
 } else {
  for (int i = 2; i <= n; i += 2) ret.push_back(i);</pre>
  for (int i = 1; i <= n; i += 2) ret.push_back(i);</pre>
8.12
       Tree Knapsack [f42766]
vector<int> G[N]; int dp[N][K]; pair<int,int> obj[N];
void dfs(int u, int mx) {
 for (int s : G[u]) {
  auto [w, v] = obj[s];
  if (mx < w) continue;</pre>
  for (int i = 0; i <= mx - w; i++)</pre>
   dp[s][i] = dp[u][i];
  dfs(s, mx - w);
  for (int i = w; i <= mx; i++)</pre>
   dp[u][i] = max(dp[u][i], dp[s][i - w] + v);
 }
       Manhattan MST [1008bc]
8.13
vector<array<int, 3>> manhattanMST(vector<P> ps) {
 vector<int> id(ps.size()); iota(all(id), 0);
 vector<array<int, 3>> edges;
 for (int k = 0; k < 4; k++) {
  sort(all(id), [&](int i, int j) {
   return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y; });</pre>
  map<int, int> sweep;
for (int i : id) {
   for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++))
    if (P d = ps[i] - ps[it->second]; d.y > d.x) break;
    else edges.push_back({d.y + d.x, i, it->second});
   }
   sweep[-ps[i].y] = i;
  for (P &p : ps)
   if (k \& 1) p.x = -p.x;
   else swap(p.x, p.y);
 return edges; // [{w, i, j}, ...]
} // test @ yosupo judge
8.14
       Binary Search On Fraction [765c5a]
struct Q {
 ll p, q;
 Q go(Q b, ll d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
 Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
 assert(pred(hi));
 bool dir = 1, L = 1, H = 1;
 for (; L || H; dir = !dir) {
  ll len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
   if (Q mid = hi.go(lo, len + step);
     mid.p > N || mid.q > N || dir ^ pred(mid))
   else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
 return dir ? hi : lo;
8.15 Barrett Reduction [d44617]
 using Big = __uint128_t; llu b, m;
 FastMod(llu b) : b(b), m(-1ULL / b) {}
 llu reduce(llu a) { // a % b
  llu r = a - (llu)((Big(m) * a) >> 64) * b;
  return r >= b ? r - b : r;
|};
```