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	Math 5.1	13 13 13 13 13 13 13 14 14 14 15 16 16 16 16 17 17 17	#ifdef k #define "" li #define "template void qwe cerr < int cr (, } template void dvo cerr < for (i	<pre>safe cerr<<pretty_function\ "<<line<<"="" .d::cerr;="" .ne="" <typenamet="" a)="" debug(a)="" dvorak(#a,="" e="" orange(a)="" qwerty(#a,="" safe\n"=""> erty(const char *s, Ta) {</pretty_function\></pre>
	Geometry 6.1 Basic Geometry 6.2 2D Convex Hull 6.3 3D Convex Hull 6.4 2D Farthest Pair 6.5 2D Closest Pair 6.6 kD Closest Pair 6.6 kD Closest Pair (3D ver.) 6.7 Simulated Annealing 6.8 Half Plane Intersection 6.9 Minkowski Sum 6.10 Intersection of line and Circle 6.11 Intersection of Polygon and Circle 6.12 Intersection of Two Circle 6.13 Tangent line of Two Circle 6.14 Minimum Covering Circle 6.15 KDTree (Nearest Point) 6.16 Rotating Sweep Line 6.17 Circle Cover	19 19 19 19 20 20 20 20 20 20	cerr < } #else #define #define #define const ir register char *pasm(// main	<pre>c << (f++ ? ", " : "") << *L; safe ((void)0) debug() ((void)0) orange() ((void)0) crease Stack at size = 256 << 20; long rsp asm("rsp"); = (char*)malloc(size)+size, *bak = (char*)rsp; "movq %0, %%rsp\n"::"r"(p)); "movq %0, %%rsp\n"::"r"(bak));</pre>

7 Stringology

7.1 Hash

1.4 Pragma Optimization

```
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
1.5 IO Optimization
static inline int gc() {
  constexpr int B = 1 << 20;
  static char buf[B], *p, *q;
  if(p == a \&\&
    (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
   return EOF;
  return *p++;
template < typename T >
static inline bool gn( T &x ) {
 int c = gc(); T sgn = 1; x = 0;
while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
if(c == '-') sgn = -1, c = gc();
 if(c == EOF) return false;
 while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
 return x *= sgn, true;
```

#pragma GCC optimize("Ofast,no-stack-protector")

#pragma GCC optimize("no-math-errno,unroll-loops")

2 Data Structure

2.1 Dark Magic

2.2 Link-Cut Tree

p->ch[dir]=c;

```
struct Node{
Node *par, *ch[2];
int xor_sum, v;
bool is_rev;
Node(int _v){
 v=xor_sum=_v;is_rev=false;
 par=ch[0]=ch[1]=nullptr;
inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
inline void down(){
 if(is_rev){
  if(ch[0]!=nullptr) ch[0]->set_rev();
   if(ch[1]!=nullptr) ch[1]->set_rev();
   is_rev=false;
 }
inline void up(){
 xor_sum=v;
  if(ch[0]!=nullptr){
  xor_sum^=ch[0]->xor_sum;
  ch[0]->par=this;
 if(ch[1]!=nullptr){
  xor_sum^=ch[1]->xor_sum;
  ch[1]->par=this;
inline bool is_root(){
 return par==nullptr ||\
   (par->ch[0]!=this && par->ch[1]!=this);
bool is_rch(){return !is_root() && par->ch[1]==this;}
} *node[maxn], *stk[maxn];
int top;
void to_child(Node* p,Node* c,bool dir){
```

```
p->up();
inline void rotate(Node* node){
 Node* par=node->par;
 Node* par_par=par->par;
 bool dir=node->is_rch()
 bool par_dir=par->is_rch()
 to_child(par, node->ch[!dir], dir);
 to_child(node,par,!dir);
 if(par_par!=nullptr && par_par->ch[par_dir]==par)
  to_child(par_par,node,par_dir);
 else node->par=par_par;
inline void splay(Node* node){
 Node* tmp=node;
 stk[top++]=node;
 while(!tmp->is_root()){
  tmp=tmp->par;
  stk[top++]=tmp;
 while(top) stk[--top]->down();
 for(Node *fa=node->par;
  !node->is_root();
  rotate(node), fa=node->par)
  if(!fa->is_root())
   rotate(fa->is_rch()==node->is_rch()?fa:node);
inline void access(Node* node){
 Node* last=nullptr;
 while(node!=nullptr){
  splay(node);
  to_child(node, last, true);
  last=node;
  node=node->par;
inline void change_root(Node* node){
 access(node);splay(node);node->set_rev();
inline void link(Node* x, Node* y){
 change_root(x);splay(x);x->par=y;
inline void split(Node* x,Node* y){
 change_root(x);access(y);splay(x);
 to_child(x,nullptr,true);y->par=nullptr;
inline void change_val(Node* node,int v){
access(node);splay(node);node->v=v;node->up();
inline int query(Node* x,Node* y){
 change_root(x);access(y);splay(y);
 return y->xor_sum;
inline Node* find_root(Node* node){
 access(node);splay(node);
 Node* last=nullptr:
 while(node!=nullptr){
  node->down();last=node;node=node->ch[0];
 return last;
set<pii> dic;
inline void add_edge(int u,int v){
 if(u>v) swap(u,v)
 if(find_root(node[u])==find_root(node[v])) return;
 dic.insert(pii(u,v))
link(node[u],node[v]);
inline void del_edge(int u,int v){
 if(u>v) swap(u,v);
 if(dic.find(pii(u,v))==dic.end()) return;
 dic.erase(pii(u,v))
 split(node[u],node[v]);
2.3 LiChao Segment Tree
struct Line{
 int m, k, id;
 Line() : id( -1 ) {}
Line('int a, int'b,'int c')
: m(a), k(b), id(c) {}
```

int at(int x) { return m * x + k; }

```
class LiChao {
private:
                                                            2.5 Sparse Table
  int n; vector< Line > nodes;
  inline int lc( int x ) { return 2 * x + 1; }
                                                            template < typename T, typename Cmp_ = less< T > >
  inline int rc( int x ) { return 2 * x + 2; }
                                                            class SparseTable {
  void insert( int 1, int r, int id, Line ln ) {
                                                            private:
   int m = (1 + r) >> 1;
                                                             vector< vector< T > > tbl;
   if ( nodes[ id ].id == -1 ) {
                                                             vector< int > lg;
   nodes[ id ] = ln;
                                                             T cv( T a, T b ) {
    return
                                                              return Cmp_()( a, b ) ? a : b;
   bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
                                                            public:
   if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
                                                             void init( T arr[], int n ) {
   atLeft ^= 1; swap( nodes[ id ], ln );
                                                              // 0-base
                                                              lg.resize(n + 1);
   if ( r - l == 1 ) return;
                                                              lg[0] = -1;
   if ( atLeft ) insert( l, m, lc( id ), ln );
                                                              for( int i=1 ; i<=n ; ++i ) lg[i] = lg[i>>1] + 1;
   else insert( m, r, rc( id ), ln );
                                                              tbl.resize(lg[n] + 1);
                                                              tbl[ 0 ].resize( n );
  int query( int 1, int r, int id, int x ) {
                                                              copy( arr, arr + n, tbl[ 0 ].begin() );
   int ret = 0;
                                                              for ( int i = 1 ; i <= lg[ n ] ; ++ i ) {
   if ( nodes[ id ].id != -1 )
                                                               int len = 1 << ( i - 1 ), sz = 1 << i;</pre>
   ret = nodes[ id ].at( x );
                                                               tbl[ i ].resize( n - sz + 1 );
   int m = (1 + r) >> 1;
                                                               for ( int j = 0 ; j \le n - sz ; ++ j
   if ( r - l == 1 ) return ret;
                                                                tbl[i][j] = cv(tbl[i-1][j], tbl[i-1][j+len]);
   else if (x < m)
                                                              }
    return max( ret, query( 1, m, lc( id ), x ) );
   else
                                                             T query( int 1, int r ) {
    return max( ret, query( m, r, rc( id ), x ) );
                                                              // 0-base [1, r)
                                                              int wh = \lg[r-1],
                                                                                    len = 1 << wh;
public:
                                                              return cv( tbl[ wh ][ 1 ], tbl[ wh ][ r - len ] );
 void build( int n_ ) {
  n = n_; nodes.clear();
                                                           };
  nodes.resize( n << 2, Line() );</pre>
                                                            2.6 Linear Basis
 void insert( Line ln ) { insert( 0, n, 0, ln ); }
                                                            template <int BITS>
  int query( int x ) { return query( 0, n, 0, x ); }
                                                            struct LinearBasis {
                                                             array<uint64_t, BITS> basis;
2.4 Treap
                                                             Basis() { basis.fill(0); }
                                                             void add(uint64_t x)
namespace Treap{
                                                              for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
  if (basis[i] == 0) {
#define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
struct node{
                                                                basis[i] = x;
 int size;
                                                                return;
  uint32_t pri;
 node *lc, *rc, *pa;
                                                               x ^= basis[i];
 node():size(0),pri(rand()),lc(0),rc(0),pa(0){}
                                                              }
 void pull()
  size = 1; pa = nullptr;
                                                             bool ok(uint64_t x) {
  if ( lc ) { size += lc->size; lc->pa = this; }
if ( rc ) { size += rc->size; rc->pa = this; }
                                                              for (int i = 0; i < BITS; ++i)</pre>
                                                               if ((x >> i) & 1) x ^= basis[i];
  }
                                                              return x == 0:
node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
                                                           };
 if ( L->pri > R->pri ) {
                                                                  Binary Search On Segment Tree
  L->rc = merge( L->rc, R ); L->pull();
   return L;
                                                            // find_first = x -> minimal x s.t. check( [a, x) )
                                                            } else {
  R->lc = merge( L, R->lc ); R->pull();
                                                            template <typename C>
                                                            int find_first(int 1, const C &check) {
   return R;
 }
                                                             if (1 >= n)
                                                              return n;
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                             1 += sz;
 if ( not rt ) L = R = nullptr;
                                                             for (int i = height; i > 0; i--)
  else if( sz( rt->lc ) + 1 <= k ) {
                                                              propagate(1 >> i)
                                                             Monoid sum = identity;
  split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                             do {
  L->pull();
                                                              while ((1 \& 1) == 0)
 } else {
                                                               1 >>= 1:
                                                              if (check(f(sum, data[1]))) {
  R = rt
   split_by_size( rt->lc, k, L, R->lc );
                                                               while (1 < sz) {</pre>
                                                                propagate(1);
   R->pull();
  }
                                                                auto nxt = f(sum, data[1]);
int getRank(node *o) {
                                                                if (not check(nxt)) {
 int r = sz(o->lc)
                                                                 sum = nxt;
  for (;o->pa != nullptr; o = o->pa)
                                                                 1++;
  if (o->pa->rc != o) r += sz(o->pa->lc);
  return r;
                                                               return 1 + 1 - sz;
```

#undef sz

```
National Taiwan University - ckiseki
  sum = f(sum, data[1++]);
 } while ((1 & -1) != 1);
 return n;
template <typename C>
int find_last(int r, const C &check) {
if (r <= 0)
 return -1;
 r += sz;
 for (int i = height; i > 0; i--)
  propagate((r - 1) >> i);
 Monoid sum = identity;
 do {
  while (r > 1 \text{ and } (r \& 1))
   r >>= 1:
  if (check(f(data[r], sum))) {
   while (r < sz) {</pre>
    propagate(r);
    r = (r << 1) + 1;
    auto nxt = f(data[r], sum);
    if (not check(nxt)) {
     sum = nxt;
     r--;
    }
   return r - sz;
  sum = f(data[r], sum);
 } while ((r & -r) != r);
 return -1;
3
     Graph
    BCC Edge
class BCC_Bridge {
 private:
  int n, ecnt;
  vector<vector<pair<int,int>>> G;
vector<int> dfn, low;
```

```
vector<bool> bridge;
 void dfs(int u, int f)
  dfn[u] = low[u] = dfn[f] + 1;
  for (auto [v, t]: G[u]) {
   if (v == f) continue;
    if (dfn[v]) {
    low[u] = min(low[u], dfn[v]);
     continue;
   dfs(v, u);
   low[u] = min(low[u], low[v]);
    if (low[v] > dfn[u]) bridge[t] = true;
public:
 void init(int n_) {
  G.clear(); G.resize(n = n_);
  low.assign(n, ecnt = \theta);
   dfn.assign(n, 0);
 void add_edge(int u, int v) {
  G[u].emplace_back(v, ecnt);
  G[v].emplace_back(u, ecnt++);
  void solve() {
  bridge.assign(ecnt, false);
  for (int i = 0; i < n; ++i)</pre>
   if (not dfn[i]) dfs(i, i);
 bool is_bridge(int x) { return bridge[x]; }
} bcc_bridge;
3.2 BCC Vertex
class BCC_AP {
```

```
class BCC_AP {
private:
  int n, ecnt;
  vector<vector<pair<int,int>>> G;
  vector<int> bcc, dfn, low, st;
```

```
vector<bool> ap, ins;
  void dfs(int u, int f) {
   dfn[u] = low[u] = dfn[f] + 1;
   int ch = 0;
   for (auto [v, t]: G[u]) if (v != f) {
    if (not ins[t]) {
     st.push_back(t);
     ins[t] = true;
    if (dfn[v]) {
     low[u] = min(low[u], dfn[v]);
     continue;
    } ++ch; dfs(v, u);
    low[u] = min(low[u], low[v]);
    if (low[v] >= dfn[u]) {
     ap[u] = true;
     while (true) {
      int eid = st.back(); st.pop_back();
      bcc[eid] = ecnt;
      if (eid == t) break;
     ecnt++;
   if (ch == 1 and u == f) ap[u] = false;
  }
 public:
  void init(int n_) {
   G.clear(); G.resize(n = n_{-});
   ecnt = 0; ap.assign(n, false);
   low.assign(n, 0); dfn.assign(n, 0);
  void add_edge(int u, int v) {
   G[u].emplace_back(v, ecnt);
   G[v].emplace_back(u, ecnt++);
  void solve() {
   ins.assign(ecnt, false);
   bcc.resize(ecnt); ecnt = 0;
   for (int i = 0; i < n; ++i)
    if (not dfn[i]) dfs(i, i);
  int get_id(int x) { return bcc[x]; }
  int count() { return ecnt; }
  bool is_ap(int x) { return ap[x]; }
} bcc_ap;
3.3 2-SAT (SCC)
class TwoSat{
private:
  int n;
  vector<vector<int>> rG,G,sccs;
  vector<int> ord,idx;
  vector<bool> vis,result;
  void dfs(int u){
   vis[u]=true
   for(int v:G[u])
   if(!vis[v]) dfs(v);
   ord.push_back(u);
  void rdfs(int u){
  vis[u]=false;idx[u]=sccs.size()-1;
   sccs.back().push_back(u);
   for(int v:rG[u])
    if(vis[v])rdfs(v);
 public:
  void init(int n_){
   n=n_;G.clear();G.resize(n);
   rG.clear();rG.resize(n);
   sccs.clear();ord.clear();
   idx.resize(n);result.resize(n);
  void add_edge(int u,int v){
   G[u].push_back(v);rG[v].push_back(u);
  void orr(int x,int y){
   if ((x^y)==1)return
   add_edge(x^1,y); add_edge(y^1,x);
```

bool solve(){

vis.clear();vis.resize(n);

```
for(int i=0;i<n;++i)</pre>
                                                                  G[ v ].push_back( u );
    if(not vis[i])dfs(i);
   reverse(ord.begin(),ord.end());
                                                                 void decompose(){
   for (int u:ord){
                                                                 chain_ = 1;
    if(!vis[u])continue;
                                                                  predfs( 1, 1 );
    sccs.push_back(vector<int>());
                                                                  time_ = 0;
    rdfs(u);
                                                                  dfschain(1,1);
   for(int i=0;i<n;i+=2)</pre>
                                                                 PII get_subtree(int u) { return {tl[ u ],tr[ u ] }; }
                                                                 vector< PII > get_path( int u , int v ){
    if(idx[i]==idx[i+1])
                                                                 vector< PII > res;
     return false;
                                                                  int g = lca( u, v );
   vector<bool> c(sccs.size());
                                                                 while ( chain[ u ] != chain[ g ] ) {
  int s = chain_st[ chain[ u ] ];
   for(size_t i=0;i<sccs.size();++i){</pre>
    for(size_t j=0;j<sccs[i].size();++j){</pre>
     result[sccs[i][j]]=c[i];
                                                                   res.emplace_back( tl[ s ], tl[ u ] + 1 );
     c[idx[sccs[i][j]^1]]=!c[i];
                                                                   u = fa[ s ][ 0 ];
                                                                  res.emplace_back( tl[ g ], tl[ u ] + 1 );
                                                                  while ( chain[ v ] != chain[ g ] ) {
   return true;
                                                                  int s = chain_st[ chain[ v ] ];
                                                                   res.emplace_back( tl[ s ], tl[ v ] + 1 );
 bool get(int x){return result[x];}
  inline int get_id(int x){return idx[x];}
                                                                   v = fa[ s ][ 0 ];
  inline int count(){return sccs.size();}
} sat2:
                                                                  res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
                                                                  return res;
3.4 Lowbit Decomposition
                                                                  /* res : list of intervals from u to v
                                                                   * ( note only nodes work, not edge )
class LowbitDecomp{
private:
                                                                   * vector< PII >& path = tree.get_path( u , v )
int time_, chain_, LOG_N;
                                                                   * for( auto [ 1, r ] : path ) {
vector< vector< int > > G, fa;
                                                                   * 0-base [ 1, r )
vector< int > tl, tr, chain, chain_st;
\//\ {\rm chain}_- : number of chain
                                                                   * }
                                                                   */
 // tl, tr[ u ] : subtree interval in the seq. of u
                                                                }
// chain_st[ u ] : head of the chain contains u
 // chian[ u ] : chain id of the chain u is on
                                                               } tree;
void predfs( int u, int f ) {
                                                                     MaxClique
                                                               3.5
 chain[ u ] = 0;
  for ( int v : G[ u ] ) {
                                                               // contain a self loop u to u, than u won't in clique
  if ( v == f ) continue;
                                                               template < size_t MAXN >
   predfs( v, u );
                                                               class MaxClique{
   if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
                                                               private:
   chain[ u ] = chain[ v ];
                                                                 using bits = bitset< MAXN >;
                                                                 bits popped, G[ MAXN ], ans
 if ( not chain[ u ] )
                                                                 size_t deg[ MAXN ], deo[ MAXN ], n;
   chain[ u ] = chain_ ++;
                                                                 void sort_by_degree() {
                                                                  popped.reset();
                                                                  for ( size_t i = 0 ; i < n ; ++ i )</pre>
void dfschain( int u, int f ) {
 fa[ u ][ 0 ] = f;
for ( int i = 1 ; i < LOG_N ; ++ i</pre>
                                                                    deg[ i ] = G[ i ].count();
                                                                  for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;</pre>
  fa[u][i] = fa[fa[u][i-1]][i-1];
                                                                    for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )
    mi = deg[ id = j ];</pre>
  tl[ u ] = time_++;
  if ( not chain_st[ chain[ u ] ] )
   chain_st[ chain[ u ] ] = u;
  for ( int v : G[ u ] )
                                                                    popped[ deo[ i ] = id ] = 1;
  if ( v != f and chain[ v ] == chain[ u ] )
                                                                    for( size_t u = G[ i ]._Find_first() ;
                                                                     u < n ; u = G[ i ]._Find_next( u ) )</pre>
    dfschain( v, u );
  for ( int v : G[ u ] )
                                                                      -- deg[ u ];
  if ( v != f and chain[ v ] != chain[ u ] )
                                                                 }
    dfschain( v, u );
  tr[ u ] = time_;
                                                                 void BK( bits R, bits P, bits X ) {
                                                                 if (R.count()+P.count() <= ans.count()) return;
if ( not P.count() and not X.count() ) {</pre>
bool anc( int u, int v )
 return tl[ u ] <= tl[ v ] and tr[ v ] <= tr[ u ];</pre>
                                                                   if ( R.count() > ans.count() ) ans = R;
                                                                   return;
public:
int lca( int u, int v ) {
                                                                  /* greedily chosse max degree as pivot
  if ( anc( u, v ) ) return u;
                                                                  bits cur = P \mid X; size_t pivot = 0, sz = 0;
  for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
                                                                  for ( size_t u = cur._Find_first() ;
  if ( not anc( fa[ u ][ i ], v ) )
                                                                   u < n ; u = cur._Find_next( u )</pre>
   u = fa[ u ][ i ];
                                                                    if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
  return fa[ u ][ 0 ];
                                                                  cur = P & ( ~G[ pivot ] );
                                                                  */ // or simply choose first
                                                                  bits cur = P & (~G[ ( P | X )._Find_first() ]);
 void init( int n ) {
 fa.assign( ++n, vector< int >( LOG_N ) );
for ( LOG_N = 0 ; ( 1 << LOG_N ) < n ; ++ LOG_N );
                                                                  for ( size_t u = cur._Find_first()
                                                                   u < n ; u = cur._Find_next( u ) ) {
 G.clear(); G.resize( n );
                                                                   if ( R[ u ] ) continue;
 tl.assign( n, 0 ); tr.assign( n, 0 ); chain.assig( n, 0 ); chain_st.assign( n, 0 );
                                                                   R[u] = 1;
                                                                   BK(R, P & G[u], X & G[u]);
                                                                   R[u] = P[u] = 0, X[u] = 1;
 void add_edge( int u , int v ) {
  // 1-base
 G[ u ].push_back( v );
                                                               public:
```

```
void init( size_t n_ ) {
 n = n_{-};
                                                               int solve(bitset<kN> mask) { // vertex mask
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                vector<int> r, c;
  G[ i ].reset();
                                                                for (int i = 0; i < n; i++)
                                                                if (mask[i]) r.push_back(i);
 ans.reset();
                                                                for (int i = 0; i < n; i++)
 void add_edges( int u, bits S ) { G[ u ] = S; }
                                                                d[i] = int((a[i] & mask).count());
void add_edge( int u, int v ) {
                                                                sort(r.begin(), r.end();
 G[u][v] = G[v][u] = 1;
                                                                [&](int i, int j) { return d[i] > d[j]; });
                                                               csort(r, c);
int solve() {
                                                               dfs(r, c, 1, mask);
 sort_by_degree(); // or simply iota( deo... )
                                                                return ans; // sol[0 ~ ans-1]
  for ( size_t i = 0 ; i < n ; ++ i )
   deg[ i ] = G[ i ].count();
                                                             } graph;
  bits pob, nob = 0; pob.set();
                                                             3.7 Virtural Tree
 for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
 for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t v = deo[ i ];</pre>
                                                             inline bool cmp(const int &i, const int &j) {
                                                              return dfn[i] < dfn[j];</pre>
  bits tmp; tmp[ v ] = 1;
  BK(tmp, pob & G[v], nob & G[v]);

pob[v] = 0, nob[v] = 1;
                                                             void build(int vectrices[], int k) {
                                                              static int stk[MAX_N];
                                                               sort(vectrices, vectrices + k, cmp);
  return static_cast< int >( ans.count() );
                                                               stk[sz++] = 0;
                                                               for (int i = 0; i < k; ++i) {
                                                                int u = vectrices[i], lca = LCA(u, stk[sz - 1]);
                                                                if (lca == stk[sz - 1]) stk[sz++] = u;
3.6 MaxCliqueDyn
                                                               else {
constexpr int kN = 150;
                                                                 while (sz \ge 2 \&\& dep[stk[sz - 2]] \ge dep[lca]) {
                                                                  addEdge(stk[sz - 2], stk[sz - 1]);
struct MaxClique { // Maximum Clique
\verb|bitset<kN>| a[kN], cs[kN]|
                                                                  sz--;
int ans, sol[kN], q, cur[kN], d[kN], n;
                                                                 if (stk[sz - 1] != lca) {
void init(int _n) {
 n = n, and q = 0;
                                                                  addEdge(lca, stk[--sz]);
 for (int i = 0; i < n; i++) a[i].reset();</pre>
                                                                  stk[sz++] = lca, vectrices[cnt++] = lca;
void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
void csort(vector<int> &r, vector<int> &c) {
                                                                 stk[sz++] = u;
                                                               }
 int mx = 1, km = max(ans - q + 1, 1), t = 0,
   m = int(r.size())
                                                               for (int i = 0; i < sz - 1; ++i)
  cs[1].reset(); cs[2].reset()
                                                               addEdge(stk[i], stk[i + 1]);
 for (int i = 0; i < m; i++) {
  int p = r[i], k = 1;
while ((cs[k] & a[p]).count()) k++;
                                                             3.8 Centroid Decomposition
  if (k > mx) cs[++mx + 1].reset();
                                                             struct Centroid {
                                                               vector<vector<int64_t>> Dist;
   cs[k][p] = 1;
                                                               vector<int> Parent, Depth;
   if (k < km) r[t++] = p;
                                                               vector<int64_t> Sub, Sub2;
                                                               vector<int> Sz, Sz2;
  c.resize(m);
  if (t) c[t - 1] = 0;
                                                               Centroid(vector<vector<pair<int, int>>> g) {
  for (int k = km; k <= mx; k++) {</pre>
                                                               int N = g.size();
  for (int p = int(cs[k]._Find_first());
                                                                vector<bool> Vis(N);
      p < kN; p = int(cs[k]._Find_next(p))) {</pre>
                                                                vector<int> sz(N), mx(N);
                                                                vector<int> Path;
    r[t] = p; c[t++] = k;
                                                                Dist.resize(N)
                                                                Parent.resize(N);
                                                                Depth.resize(N)
 void dfs(vector<int> &r, vector<int> &c, int 1,
                                                                auto DfsSz = [&](auto dfs, int x) -> void {
 bitset<kN> mask) {
                                                                 Vis[x] = true; sz[x] = 1; mx[x] = 0;
                                                                 for (auto [u, w] : g[x]) {
  while (!r.empty()) {
  int p = r.back(); r.pop_back();
                                                                  if (Vis[u]) continue;
                                                                  dfs(dfs, u)
   mask[p] = 0;
   if (q + c.back() <= ans) return;</pre>
                                                                  sz[x] += sz[u];
   cur[q++] = p;
                                                                  mx[x] = max(mx[x], sz[u]);
   vector<int> nr, nc;
   bitset<kN> nmask = mask & a[p];
                                                                Path.push_back(x);
   for (int i : r)
                                                                }:
    if (a[p][i]) nr.push_back(i);
                                                                auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
   if (!nr.empty()) {
    if (1 < 4) {
                                                                 Dist[x].push_back(D);Vis[x] = true;
     for (int i : nr)
                                                                 for (auto [u, w] : g[x]) {
      d[i] = int((a[i] & nmask).count());
                                                                  if (Vis[u]) continue;
     sort(nr.begin(), nr.end(),
                                                                  dfs(dfs, u, D + w);
      [&](int x, int y)
       return d[x] > d[y];
                                                               };
                                                                auto Dfs = [&]
      });
                                                                 (auto dfs, int x, int D = 0, int p = -1)->void {
                                                                 Path.clear(); DfsSz(DfsSz, x);
   csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
                                                                 int M = Path.size();
   ans = q; copy(cur, cur + q, sol);
                                                                 int C = -1;
                                                                 for (int u : Path) {
   c.pop_back(); q--;
                                                                  if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
                                                                  Vis[u] = false;
```

for(int i=0; i<n; i++) {</pre>

```
double avg=-inf;
   DfsDist(DfsDist, C);
                                                                  for(int k=0; k<n; k++) {</pre>
   for (int u : Path) Vis[u] = false;
                                                                   if(d[n][i]<inf-eps)</pre>
   Parent[C] = p; Vis[C] = true;
                                                                    avg=max(avg,(d[n][i]-d[k][i])/(n-k));
   Depth[C] = D;
                                                                   else avg=max(avg,inf);
   for (auto [u, w] : g[C]) {
    if (Vis[u]) continue;
                                                                  if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
    dfs(dfs, u, D + 1, C);
                                                                 FZ(vst);edgeID.clear();cycle.clear();rho.clear();
                                                                 for (int i=n; !vst[st]; st=prv[i--][st]) {
 Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
                                                                  vst[st]++
  Sz.resize(N); Sz2.resize(N);
                                                                  edgeID.PB(prve[i][st]);
                                                                  rho.PB(st);
 void Mark(int v) {
 int x = v, z = -1;
                                                                 while (vst[st] != 2) {
 for (int i = Depth[v]; i >= 0; --i) {
                                                                  int v = rho.back(); rho.pop_back();
  Sub[x] += Dist[v][i]; Sz[x]++;
                                                                  cycle.PB(v);
  if (z != -1) {
                                                                  vst[v]++;
   Sub2[z] += Dist[v][i];
                                                                 }
    Sz2[z]++;
                                                                 reverse(ALL(edgeID));
                                                                 edgeID.resize(SZ(cycle));
   z = x; x = Parent[x];
                                                                 return mmc;
                                                               } mmc;
int64_t Query(int v) {
 int64_t res = 0;
                                                               3.11 Mo's Algorithm on Tree
 int x = v, z = -1;
                                                               int q; vector< int > G[N];
 for (int i = Depth[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
                                                               struct Que{
                                                                int u, v, id;
  if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
                                                               } que[ N ];
  z = x; x = Parent[x];
                                                               int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
 }
                                                               void_dfs( int u, int f ) {
  return res;
                                                                dfn[ u ] = dfn_++; int saved_rbp = stk_;
                                                                for ( int v : G[ u ] ) {
  if ( v == f ) continue;
};
3.9
     Tree Hashing
                                                                 dfs( v, u );
                                                                 if ( stk_
                                                                             saved_rbp < SQRT_N ) continue;</pre>
uint64_t hsah(int u, int f) {
                                                                 for ( ++ block_ ; stk_ != saved_rbp ; )
uint64_t r = 127;
for (int v : G[ u ]) if (v != f) {
                                                                   block_id[ stk[ -- stk_ ] ] = block_;
 uint64_t hh = hsah(v, u);
                                                                stk[ stk_ ++ ] = u;
 r=(r+(hh*hh)%1010101333)%1011820613;
                                                               bool inPath[ N ];
return r;
                                                               void Diff( int u ) {
                                                                if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
3.10 Minimum Mean Cycle
                                                                else { /*add this edge*/ }
/* minimum mean cycle O(VE) */
                                                               void traverse( int& origin_u, int u ) {
struct MMC{
                                                                for ( int g = lca( origin_u, u ) ;
#define FZ(n) memset((n),0,sizeof(n))
                                                                 origin_u != g ; origin_u = parent_of[ origin_u ] )
#define E 101010
                                                                  Diff( origin_u );
#define V 1021
                                                                for (int v = u; v != origin_u; v = parent_of[v])
#define inf 1e9
                                                                 Diff( v );
struct Edge { int v,u; double c; };
                                                                origin_u = u;
 int n, m, prv[V][V], prve[V][V], vst[V];
Edge e[E];
                                                               void solve() {
vector<int> edgeID, cycle, rho;
                                                                dfs( 1, 1 );
double d[V][V];
                                                                while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
void init( int _n ) { n = _n; m = 0; }
                                                                sort( que, que + q, [](const Que& x, const Que& y) {
  return tie( block_id[ x.u ], dfn[ x.v ] )
 // WARNING: TYPE matters
void add_edge( int vi , int ui , double ci )
{ e[ m ++ ] = { vi , ui , ci }; }
                                                                      < tie( block_id[ y.u ], dfn[ y.v ] );
                                                                } );
void bellman_ford() {
                                                                int U = 1, V = 1;
 for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {</pre>
                                                                for ( int i = 0 ; i < q ; ++ i ) {
   fill(d[i+1], d[i+1]+n, inf);
                                                                 pass( U, que[ i ].u );
  for(int j=0; j<m; j++) {
  int v = e[j].v, u = e[j].u;
  if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
                                                                 pass( V, que[ i ].v );
                                                                 // we could get our answer of que[ i ].id
                                                                }
                                                               }
     d[i+1][u] = d[i][v]+e[j].c;
     prv[i+1][u] = v;
                                                               Method 2:
     prve[i+1][u] = j;
                                                               dfs u:
                                                                push u
   }
                                                                iterate subtree
  }
                                                                push u
                                                               Let P = LCA(u, v), and St(u) \le St(v)
double solve(){
                                                               if (P == u) query[St(u), St(v)]
  // returns inf if no cycle, mmc otherwise
                                                               else query[Ed(u), St(v)], query[St(P), St(P)]
 double mmc=inf;
  int st = -1
 bellman_ford();
```

3.12 Minimum Steiner Tree

in[e.v] = e.w;

```
// Minimum Steiner Tree
                                                                          prv[e.v] = e.u;
// 0(V 3^T + V^2 2^T)
struct SteinerTree{
                                                                        in[root] = 0;
                                                                        prv[root] = -1;
for (int i = 0; i < n; i++)
#define V 33
#define T 8
#define INF 1023456789
                                                                         if (in[i] == -inf)
int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
                                                                          return -inf;
                                                                         // find cycle
void init( int _n ){
                                                                        int tot = 0;
  for( int i = 0 ; i < n ; i ++ ){
  for( int j = 0 ; j < n ; j ++ )</pre>
                                                                        vector<int> id(n, -1), vis(n, -1);
for (int i = 0; i < n; i++) {</pre>
    dst[ i ][ j ] = INF;
                                                                         ans += in[i];
                                                                         for (int x = i; x != -1 && id[x] == -1; x = prv[x])
   dst[ i ][ i ] = 0;
                                                                           if (vis[x] == i) {
void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
  dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
                                                                            for (int y = prv[x]; y != x; y = prv[y])
                                                                             id[y] = tot;
                                                                            id[x] = tot++;
                                                                            break;
 void shortest_path(){
 for( int k = 0 ; k < n ; k ++ )
                                                                           vis[x] = i;
   for( int i = 0 ; i < n ; i ++ )
for( int j = 0 ; j < n ; j ++ )
dst[ i ][ j ] = min( dst[ i ][ j ],
                                                                         }
                                                                        if (!tot)
                                                                         return ans;
         dst[ i ][ k ] + dst[ k ][ j ] );
                                                                        for (int i = 0; i < n; i++)
if (id[i] == -1)
int solve( const vector<int>& ter ){
  int t = (int)ter.size();
                                                                          id[i] = tot++;
  for( int i = 0 ; i < (1 << t ) ; i ++ )
for( int j = 0 ; j < n ; j ++ )
dp[ i ][ j ] = INF;</pre>
                                                                         // shrink
                                                                        for (auto &e : E) {
                                                                         if (id[e.u] != id[e.v])
  for( int i = 0 ; i < n ; i ++ )
dp[ 0 ][ i ] = 0;
for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
                                                                          e.w -= in[e.v];
                                                                         e.u = id[e.u], e.v = id[e.v];
                                                                        n = tot;
   if( msk == ( msk & (-msk) ) ){
    int who = __lg( msk );
for( int i = 0 ; i < n ; i ++ )</pre>
                                                                        root = id[root];
     dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
                                                                       assert(false);
    continue:
                                                                    } DMST;
   for( int i = 0 ; i < n ; i ++ )</pre>
                                                                            Manhattan Minimum Spanning Tree
    for( int submsk = ( msk - 1 ) & msk ; submsk ;
          submsk = ( submsk - 1 ) & msk )
                                                                    typedef Point<int> P;
       dp[ msk ][ i ] = min( dp[ msk ][ i ],
                                                                     vector<array<int, 3>> manhattanMST(vector<P> ps) {
                dp[ submsk ][ i ] +
                                                                      vi id(sz(ps));
                dp[ msk ^ submsk ][ i ] );
                                                                      iota(all(id), 0);
   for( int i = 0 ; i < n ; i ++ ){</pre>
                                                                      vector<array<int, 3>> edges;
    tdst[ i ] = INF;
for( int j = 0 ; j < n ; j ++ )
tdst[ i ] = min( tdst[ i ],</pre>
                                                                      rep(k, 0, 4) {
                                                                       sort(all(id), [&](int i, int j) {
                                                                        return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
            dp[ msk ][ j ] + dst[ j ][ i ] );
                                                                       });
                                                                       map<int, int> sweep;
                                                                       for (int i : id) {
   for( int i = 0 ; i < n ; i ++ )</pre>
    dp[ msk ][ i ] = tdst[ i ];
                                                                        for (auto it = sweep.lower_bound(-ps[i].y);
                                                                            it != sweep.end(); sweep.erase(it++)) {
                                                                         int j = it->second;
  int ans = INF:
                                                                         P d = ps[i] - ps[j];
  for( int i = 0 ; i < n ; i ++ )</pre>
   ans = min(ans, dp[(1 << t) - 1][i]);
                                                                         if (d.y > d.x) break;
  return ans;
                                                                         edges.push_back(\{d.y + d.x, i, j\});
} solver;
                                                                        sweep[-ps[i].y] = i;
3.13 Directed Minimum Spanning Tree
                                                                       for (P &p : ps)
struct DirectedMST { // find maximum
                                                                        if (k \& 1) p.x = -p.x;
struct Edge {
                                                                        else swap(p.x, p.y);
  int u, v;
  int w;
                                                                      return edges; // [{w, i, j}, ...]
  Edge(int u, int v, int w) : u(u), v(v), w(w) {}
                                                                     3.15
                                                                            Dominator Tree
vector<Edge> Edges;
void clear() { Edges.clear(); }
                                                                    namespace dominator {
void addEdge(int a, int b, int w) { Edges.emplace_back
                                                                    vector<int> g[maxn], r[maxn], rdom[maxn];
    (a, b, w); }
                                                                     int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
 int solve(int root, int n) {
                                                                     int dom[maxn], val[maxn], rp[maxn], tk;
  vector<Edge> E = Edges;
                                                                     void init(int n) {
                                                                      // vertices are numbered from 0 to n-1
  int ans = 0:
                                                                      fill(dfn, dfn + n, -1);fill(rev, rev + n, -1);
fill(fa, fa + n, -1);fill(val, val + n, -1);
  while (true) {
   // find best in edge
   vector<int>in(n, -inf), prv(n, -1);
                                                                      fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
                                                                      fill(dom, dom + n, -1); tk = 0; for (int i = 0; i < n; ++i) {
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
```

g[i].clear(); r[i].clear(); rdom[i].clear();

color(u, L[a].first, L[a].second);

```
void add_edge(int x, int y) { g[x].push_back(y); }
                                                                 if (!G[u][v0]) {
void dfs(int x) {
rev[dfn[x] = tk] = x;
 fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
 for (int u : g[x]) {
 if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
  r[dfn[u]].push_back(dfn[x]);
                                                                  } else t--;
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
  if (fa[x] == x) return c ? -1 : x;
                                                               4
                                                                     Matching & Flow
 int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
                                                               4.1
                                                                     Kuhn Munkres
                                                               class KM {
 fa[x] = p;
                                                               private:
 return c ? p : val[x];
vector<int> build(int s, int n) {
                                                                vector<lld> hl,hr,slk;
                                                                vector<int> fl,fr,pre,qu;
// return the father of each node in the dominator tree
                                                                vector<vector<lld>> w;
// p[i] = -2 if i is unreachable from s
                                                                vector<bool> v1, vr;
 dfs(s);
 for (int i = tk - 1; i >= 0; --i)
                                                                int n, ql, qr;
  for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
                                                                bool check(int x) {
  if (i) rdom[sdom[i]].push_back(i);
  for (int &u : rdom[i]) {
   int p = find(u)
                                                                 return false:
   if (sdom[p] == i) dom[u] = i;
   else dom[u] = p;
                                                                void bfs(int s) {
  if (i) merge(i, rp[i]);
 }
 vector<int> p(n, -2); p[s] = -1;
 for (int i = 1; i < tk; ++i)
                                                                 ql = qr = 0;
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
                                                                 qu[qr++] = s;
                                                                 vr[s] = true;
 for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
 return p;
                                                                 while (true) {
                                                                  11d d;
3.16 Edge Coloring
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
 for (int i = 0; i <= N; i++)
  for (int j = 0; j \leftarrow N; j++)
    C[i][j] = G[i][j] = 0;
                                                                  d = INF:
                                                                  for (int x = 0; x < n; ++x)
void solve(vector<pair<int, int>> &E, int N) {
int X[kN] = {}, a;
auto update = [&](int u)
                                                                   if (v1[x]) h1[x] += d;
  for (X[u] = 1; C[u][X[u]]; X[u]++);
                                                                   else slk[x] -= d;
 auto color = [&](int u, int v, int c) {
  int p = G[u][v];
                                                                   if (vr[x]) hr[x] -= d;
  G[u][v] = G[v][u] = c;
  C[u][c] = v, C[v][c] = u;
  C[u][p] = C[v][p] = 0;
  if (p) X[u] = X[v] = p;
  else update(u), update(v);
                                                               public:
                                                                void init( int n_ ) {
  return p;
                                                                 n = n_; qu.resize(n);
 };
 auto flip = [&](int u, int c1, int c2) {
 int p = C[u][c1];
  swap(C[u][c1], C[u][c2]);
  if (p) G[u][p] = G[p][u] = c2;
if (!C[u][c1]) X[u] = c1;
  if (!C[u][c2]) X[u] = c2;
                                                                 vl.resize(n); vr.resize(n);
  return p;
 }:
                                                                1ld solve() {
 for (int i = 1; i <= N; i++) X[i] = 1;
                                                                 for (int i = 0; i < n; ++i)
 for (int t = 0; t < E.size(); t++) {</pre>
  auto [u, v] = E[t];
  int v0 = v, c = X[u], c0 = c, d;
                                                                 11d res = 0;
  vector<pair<int, int>> L; int vst[kN] = {};
  while (!G[u][v0]) {
                                                                 return res;
   L.emplace_back(v, d = X[v]);
   if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
   c = color(u, L[a].first, c);
else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
                                                               } km;
```

```
else if (vst[d]) break;
else vst[d] = 1, v = C[u][d];
for (; v; v = flip(v, c, d), swap(c, d));
if (C[u][c0]) { a = int(L.size()) - 1;
 while (--a >= 0 && L[a].second != c);
 for(;a>=0;a--)color(u,L[a].first,L[a].second);
```

```
static constexpr 11d INF = 1LL << 60;</pre>
 if (vl[x] = true, fl[x] != -1)
  return vr[qu[qr++] = f1[x]] = true;
 while (x != -1) swap(x, fr[fl[x] = pre[x]]);
 fill(slk.begin(), slk.end(), INF);
 fill(vl.begin(), vl.end(), false);
 fill(vr.begin(), vr.end(), false);
  while (ql < qr) {
  for (int x = 0, y = qu[ql++]; x < n; ++x) {
    if(!v1[x]&&s1k[x]>=(d=h1[x]+hr[y]-w[x][y])){
     if (pre[x] = y, d) slk[x] = d;
      else if (!check(x)) return;
   if (!vl[x] \&\& d > slk[x]) d = slk[x];
  for (int x = 0; x < n; ++x) {
  for (int x = 0; x < n; ++x)
if (!v1[x] && !slk[x] && !check(x)) return;</pre>
fl.clear(); fl.resize(n, -1);
fr.clear(); fr.resize(n, -1);
hr.clear(); hr.resize(n); hl.resize(n);
w.clear(); w.resize(n, vector<lld>(n));
slk.resize(n); pre.resize(n);
void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 hl[i] = *max_element(w[i].begin(), w[i].end());
 for (int i = 0; i < n; ++i) bfs(i);</pre>
 for (int i = 0; i < n; ++i) res += w[i][f1[i]];</pre>
```

4.2 Bipartite Matching

```
class BipartiteMatching{
                                                                 if (s[u] == -1)
                                                                  pre[u] = x, s[u] = 1;
private:
 vector<int> X[N], Y[N];
                                                                  if (match[u] == n) {
 int fX[N], fY[N], n;
                                                                   for (int a = u, b = x, last; b != n; a = last, b =
 bitset<N> walked;
                                                                   pre[a])
 bool dfs(int x){
                                                                    last = match[b], match[b] = a, match[a] = b;
  for(auto i:X[x]){
                                                                   return true;
   if(walked[i])continue;
                                                                  q.push(match[u]);
   walked[i]=1;
   if(fY[i]==-1||dfs(fY[i])){
                                                                  s[match[u]] = 0;
                                                                 } else if (!s[u] && Find(u) != Find(x)) {
    fY[i]=x;fX[x]=i;
                                                                  int 1 = LCA(u, x, n);
Blossom(x, u, 1);
    return 1:
   }
  }
                                                                  Blossom(u, x, 1);
  return 0;
public:
 void init(int _n){
                                                               return false;
  n=_n; walked.reset();
  for(int i=0;i<n;i++){</pre>
                                                              int Solve(int n) {
   X[i].clear();Y[i].clear();
                                                               int res = 0;
   fX[i]=fY[i]=-1;
                                                               for (int x = 0; x < n; ++x) {
                                                                if (match[x] == n) res += Bfs(x, n);
 void add_edge(int x, int y){
                                                               return res;
 X[x].push_back(y); Y[y].push_back(y);
                                                              }}
                                                              4.4 Minimum Weight Matching (Clique version)
 int solve(){
                                                              struct Graph {
  int cnt = 0;
                                                               // 0-base (Perfect Match)
  for(int i=0;i<n;i++){</pre>
  walked.reset();
                                                               int n, edge[MXN][MXN];
   if(dfs(i)) cnt++;
                                                               int match[MXN], dis[MXN], onstk[MXN];
                                                               vector<int> stk;
                                                               void init(int _n) {
  // return how many pair matched
  return cnt;
                                                                n = _n;
                                                                for (int i=0; i<n; i++)</pre>
};
                                                                 for (int j=0; j<n; j++)</pre>
                                                                  edge[i][j] = 0;
      General Graph Matching
                                                               void set_edge(int u, int v, int w) {
namespace matching .
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
                                                                edge[u][v] = edge[v][u] = w;
vector<int> g[kN];
queue<int> q;
                                                               bool SPFA(int u){
void Init(int n) {
                                                                if (onstk[u]) return true;
for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
                                                                stk.PB(u);
                                                                onstk[u] = 1;
                                                                for (int v=0; v<n; v++){</pre>
void AddEdge(int u, int v) {
                                                                 if (u != v && match[u] != v && !onstk[v]){
                                                                  int m = match[v]
 g[u].push_back(v);
 g[v].push_back(u);
                                                                  if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
                                                                   dis[m] = dis[u] - edge[v][m] + edge[u][v];
int Find(int u) {
                                                                   onstk[v] = 1;
return u == fa[u] ? u : fa[u] = Find(fa[u]);
                                                                   stk.PB(v);
                                                                   if (SPFA(m)) return true;
int LCA(int x, int y, int n) {
  static int tk = 0; tk++;
                                                                   stk.pop_back();
                                                                   onstk[v] = 0;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
                                                                 }
  if (x != n) {
   if (v[x] == tk) return x;
                                                                onstk[u] = 0;
   v[x] = tk;
                                                                stk.pop_back();
   x = Find(pre[match[x]]);
                                                                return false;
  }
                                                               int solve() {
void Blossom(int x, int y, int 1) {
                                                                // find a match
 while (Find(x) != 1) {
                                                                for (int i=0; i<n; i+=2){
 pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
                                                                 match[i] = i+1;
                                                                 match[i+1] = i;
  if (fa[x] == x) fa[x] = 1;
  if (fa[y] == y) fa[y] = 1;
                                                                while (true){
                                                                 int found = 0;
  x = pre[y];
                                                                 for (int i=0; i<n; i++)</pre>
                                                                  dis[i] = onstk[i] = 0;
                                                                 for (int i=0; i<n; i++){
bool Bfs(int r, int n) {
                                                                  stk.clear()
 for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;
 while (!q.empty()) q.pop();
                                                                  if (!onstk[i] && SPFA(i)){
 q.push(r);
                                                                   found = 1:
 s[r] = 0;
                                                                   while (SZ(stk)>=2){
 while (!q.empty()) {
                                                                    int u = stk.back(); stk.pop_back();
  int x = q.front(); q.pop();
                                                                    int v = stk.back(); stk.pop_back();
  for (int u : g[x]) {
                                                                    match[u] = v;
```

```
match[v] = u;
    }
   if (!found) break;
  int ret = 0:
  for (int i=0; i<n; i++)</pre>
   ret += edge[i][match[i]];
  return ret>>1;
} graph;
```

4.5 Minimum Cost Circulation

```
struct Edge { int to, cap, rev, cost; };
vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
bool mark[kN];
int NegativeCycle(int n) {
memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
int upd = -1;
for (int i = 0; i <= n; ++i) {
  for (int j = 0; j < n; ++j) {
   int idx = 0;
   for (auto &e : g[j]) {
    if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
     dist[e.to] = dist[j] + e.cost;
     pv[e.to] = j, ed[e.to] = idx;
     if (i == n) {
      upd = j;
      while(!mark[upd])mark[upd]=1,upd=pv[upd];
      return upd:
    idx++;
  }
return -1;
int Solve(int n) {
int rt = -1, ans = 0;
while ((rt = NegativeCycle(n)) >= 0) {
 memset(mark, false, sizeof(mark));
  vector<pair<int, int>> cyc;
  while (!mark[rt]) {
  cyc.emplace_back(pv[rt], ed[rt]);
  mark[rt] = true;
  rt = pv[rt];
  }
  reverse(cyc.begin(), cyc.end());
  int cap = kInf;
  for (auto &i : cyc) {
   auto &e = g[i.first][i.second];
   cap = min(cap, e.cap);
  for (auto &i : cyc) {
  auto &e = g[i.first][i.second];
   e.cap -= cap;
   g[e.to][e.rev].cap += cap;
   ans += e.cost * cap;
return ans;
```

4.6 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}$.
 - 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t\to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T.If $f
 eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the
 - maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect t o s with capacity ∞ and let the flow from S to T be f'. If $f+f'
 eq \sum_{v\in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.

- 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite araph(X,Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect x o y with (cost,cap) = (c,1) if
 - c>0, otherwise connect $y\to x$ with (cost, cap)=(-c,1) 3. For each edge with c<0, sum these cost as K, then increase d(y)by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v)>0, connect S o v with (cost, cap)=0
 - (0,d(v))5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =(0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let ${\cal K}$ be the sum of all weights

 - 3. Connect source $s \to v, v \in G$ with capacity K 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity
 - 5. For $v \in {\it G}$, connect it with sink $v \to t$ with capacity K+2T- $\left(\sum_{e \in E(v)} w(e)\right) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v
 - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_i$
 - 2. Create edge (u, v) with capacity w with w being the cost of choosing \boldsymbol{u} without choosing $\boldsymbol{v}.$
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with capacity c_y . 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

Dinic

```
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
    int to, rev;
    Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
       int u = bfs.front(); bfs.pop();
       for (auto e: G[u]) {
         if (e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
         bfs.push(e.to); lv[e.to] = lv[u] + 1;
    return lv[ed] != -1;
  Cap DFS(int u, Cap f){
  if (u == ed) return f;
    Cap ret = 0;
     for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
       auto &e = G[u][i];
       if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
       Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
```

```
if (f == 0) return ret;
                                                                 G[st].emplace_back(ed,SZ(G[ed]),c,w);
                                                                G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
    if (ret == 0) lv[u] = -1;
                                                                PCW solve(){
    return ret;
                                                                Cap cc=0; Wei ww=0;
public:
                                                                 while(true){
  void init(int n_) { G.assign(n = n_, vector<E>()); }
                                                                 PCW ret=SPFA();
  if(ret.first==-1) break;
                                                                  cc+=ret.first;
                                                                  ww+=ret.first * ret.second;
    G[v].push_back({u, int(G[u].size())-1, 0});
 Cap max_flow(int st_, int ed_){
  st = st_, ed = ed_; Cap ret = 0;
                                                                 return {cc,ww};
    while (BFS()) {
                                                              } mcmf;
      idx.assign(n, 0);
                                                               4.9 GomoryHu Tree
      Cap f = DFS(st, numeric_limits<Cap>::max());
      ret += f;
                                                              int g[maxn];
      if (f == 0) break;
                                                              vector<edge> GomoryHu(int n){
                                                                vector<edge> rt;
    return ret;
                                                                for(int i=1;i<=n;++i)g[i]=1;</pre>
                                                                for(int i=2;i<=n;++i){</pre>
};
                                                                 int t=g[i];
                                                                 flow.reset(); // clear flows on all edge
4.8 Minimum Cost Maximum Flow
                                                                 rt.push_back({i,t,flow(i,t)});
class MiniCostMaxiFlow{
                                                                 flow.walk(i); // bfs points that connected to i (use
 using Cap = int; using Wei = int64_t;
                                                                   edges not fully flow)
 using PCW = pair<Cap,Wei>;
                                                                 for(int j=i+1;j<=n;++j){</pre>
 static constexpr Cap INF_CAP = 1 << 30;</pre>
                                                                 if(g[j]==t && flow.connect(j))g[j]=i; // check if i
 static constexpr Wei INF_WEI = 1LL<<60;</pre>
                                                                   can reach j
private:
 struct Edge{
  int to, back;
                                                                return rt;
  Cap cap; Wei wei;
                                                              }
  Edge() {}
  Edge(int a,int b, Cap c, Wei d):
                                                               4.10 Global Min-Cut
   to(a),back(b),cap(c),wei(d)
                                                              const int maxn = 500 + 5;
  {}
                                                              int w[maxn][maxn], g[maxn];
 };
                                                              bool v[maxn], del[maxn];
 int ori, edd;
                                                              void add_edge(int x, int y, int c) {
 vector<vector<Edge>> G;
                                                               w[x][y] += c; w[y][x] += c;
 vector<int> fa, wh;
 vector<bool> inq;
                                                              pair<int, int> phase(int n) {
 vector<Wei> dis;
                                                               memset(v, false, sizeof(v));
 PCW SPFA(){
                                                               memset(g, 0, sizeof(g));
  fill(inq.begin(),inq.end(),false)
                                                               int s = -1, t = -1;
  fill(dis.begin(), dis.end(), INF_WEI);
                                                                while (true) {
  queue<int> qq; qq.push(ori);
                                                                 int c = -1;
  dis[ori] = 0;
                                                                 for (int i = 0; i < n; ++i) {</pre>
  while(not qq.empty()){
                                                                  if (del[i] || v[i]) continue;
   int u=qq.front();qq.pop();
                                                                  if (c == -1 \mid | g[i] > g[c]) c = i;
   inq[u] = false;
   for(int i=0;i<SZ(G[u]);++i){</pre>
                                                                if (c == -1) break;
    Edge e=G[u][i];
                                                                v[s = t, t = c] = true;
    int v=e.to; Wei d=e.wei;
                                                                 for (int i = 0; i < n; ++i) {
    if(e.cap<=0||dis[v]<=dis[u]+d)</pre>
                                                                 if (del[i] || v[i]) continue;
     continue
                                                                  g[i] += w[c][i];
    dis[v] = dis[u] + d;
    fa[v] = u, wh[v] = i;
if (inq[v]) continue;
                                                                return make_pair(s, t);
    qq.push(v);
    inq[v] = true;
                                                               int mincut(int n) {
                                                                int cut = 1e9;
                                                               memset(del, false, sizeof(del));
  if(dis[edd]==INF_WEI) return {-1, -1};
                                                                for (int i = 0; i < n - 1; ++i) {
  Cap mw=INF_CAP;
                                                                int s, t; tie(s, t) = phase(n);
del[t] = true; cut = min(cut, g[t]);
for (int j = 0; j < n; ++j) {</pre>
  for(int i=edd;i!=ori;i=fa[i])
   mw=min(mw,G[fa[i]][wh[i]].cap);
  for (int i=edd;i!=ori;i=fa[i]){
                                                                  w[s][j] += w[t][j]; w[j][s] += w[j][t];
   auto &eg=G[fa[i]][wh[i]];
                                                                 }
   eg.cap -= mw;
   G[eg.to][eg.back].cap+=mw;
                                                                return cut;
  return {mw, dis[edd]};
 }
                                                               5
                                                                    Math
public:
 void init(int a,int b,int n){
                                                               5.1 \left| \frac{n}{s} \right| Enumeration
  ori=a,edd=b;
                                                              T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor
  G.clear();G.resize(n);
 fa.resize(n);wh.resize(n);
                                                               5.2 ax+by=gcd
  inq.resize(n); dis.resize(n);
                                                              // ax+ny = 1, ax+ny == ax == 1 \pmod{n}
 void add_edge(int st, int ed, Cap c, Wei w){
                                                              void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
```

```
if (y == 0) g=x,a=1,b=0;
else exgcd(y, x\%y, g, b, a), b=(x/y)*a;
5.3
    Pollard Rho
// does not work when n is prime
// return any non-trivial factor
llu pollard_rho(llu n){
static auto f=[](llu x,llu k,llu m){
  return add(k,mul(x,x,m),m);
if (!(n&1)) return 2;
mt19937 rnd(120821011);
while(true){
 llu y=2, yy=y, x=rnd()%n, t=1;
  for(llu sz=2;t==1;sz<<=1) {</pre>
   for(llu i=0;i<sz;++i){</pre>
   if(t!=1)break;
   yy=f(yy,x,n);
    t=gcd(yy>y?yy-y:y-yy,n);
  if(t!=1&&t!=n) return t;
5.4 Pi Count (Linear Sieve)
static constexpr int N = 1000000 + 5;
lld pi[N];
vector<int> primes;
bool sieved[N];
1ld cube_root(lld x){
lld s=cbrt(x-static_cast<long double>(0.1));
while(s*s*s <= x) ++s;</pre>
return s-1;
11d square_root(11d x){
lld s=sqrt(x-static_cast<long double>(0.1));
while(s*s <= x) ++s;</pre>
return s-1;
void init(){
primes.reserve(N);
primes.push_back(1);
for(int i=2;i<N;i++) {</pre>
 if(!sieved[i]) primes.push_back(i);
 pi[i] = !sieved[i] + pi[i-1];
  for(int p: primes) if(p > 1) {
  if(p * i >= N) break;
   sieved[p * i] = true;
   if(p % i == 0) break;
 }
}
11d phi(11d m, 11d n) {
static constexpr int MM = 80000, NN = 500;
static lld val[MM][NN];
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
if(n == 0) return m;
 if(primes[n] >= m) return 1;
1ld ret = phi(m,n-1)-phi(m/primes[n],n-1);
 if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
return ret;
11d pi_count(11d);
11d P2(11d m, 11d n) {
11d sm = square_root(m), ret = 0;
for(lld i = n+1;primes[i]<=sm;i++)</pre>
 ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
 return ret;
11d pi_count(11d m) {
if(m < N) return pi[m];</pre>
11d n = pi_count(cube_root(m));
return phi(m, n) + n - 1 - P2(m, n);
5.5 Strling Number
5.5.1 First Kind
```

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1)\dots(x+n-1) = \sum_{k=0}^n S_1(n,k)x^k$$

$$g(x) = x(x+1)\dots(x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)!a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.5.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

5.6 Range Sieve

5.7 Miller Rabin

```
bool isprime(llu x){
 static llu magic[]={2,325,9375,28178,\
          450775,9780504,1795265022};
 static auto witn=[](llu a,llu u,llu n,int t)
  if (!(a = mpow(a%n,u,n)))return 0;
  while(t--){
   11u a2=mul(a,a,n);
   if(a2==1 && a!=1 && a!=n-1)
    return 1:
   a = a2:
  return a!=1;
 if(x<2)return 0;</pre>
 if(!(x&1))return x==2;
 llu x1=x-1; int t=0;
 while(!(x1&1))x1>>=1,t++;
 for(llu m:magic)if(witn(m,x1,x,t))return 0;
 return 1;
```

5.8 Extended Euler

$$a^b \equiv \begin{cases} a^b \mod \varphi(m) + \varphi(m) & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m$$

5.9 Gauss Elimination

```
void gauss(vector<vector<double>> &d) {
  int n = d.size(), m = d[0].size();
  for (int i = 0; i < m; ++i) {
    int p = -1;
    for (int j = i; j < n; ++j) {
      if (fabs(d[j][i]) < eps) continue;
      if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
    }
}
```

```
if (p == -1) continue;
                                                               VL c(sz);
                                                               for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
  for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);</pre>
  for (int j = 0; j < n; ++j) {
                                                               return c:
  if (i == j) continue;
double z = d[j][i] / d[i][i];
                                                              VI convolution_mod(const VI &a, const VI &b, int p) {
   for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
                                                               int sz = 1;
                                                               while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
                                                               vector<cplx> fa(sz), fb(sz);
                                                               for (int i = 0; i < (int)a.size(); ++i)</pre>
                                                                fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
      Fast Fourier Transform
5.10
                                                               for (int i = 0; i < (int)b.size(); ++i)</pre>
const int mod = 1000000007;
                                                                fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
                                                               fft(fa, sz), fft(fb, sz);
const int M1 = 985661441; // G = 3
                                                               double r = 0.25 / sz;
                                                               cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
const int M2 = 998244353;
                                                               for (int i = 0; i \le (sz >> 1); ++i) {
const int M3 = 1004535809;
                                                                int j = (sz - i) & (sz - 1);
                                                                cplx a1 = (fa[i] + fa[j].conj());
int superBigCRT(int64_t A, int64_t B, int64_t C) {
                                                                cplx a2 = (fa[i] - fa[j].conj()) * r2;
 static_assert (M1 <= M2 && M2 <= M3);
 constexpr int64_t r12 = modpow(M1, M2-2, M2);
                                                                cplx b1 = (fb[i] + fb[j].conj()) * r3;
                                                                cplx b2 = (fb[i] - fb[j].conj()) * r4;
 constexpr int64_t r13 = modpow(M1, M3-2, M3);
 constexpr int64_t r23 = modpow(M2, M3-2, M3);
                                                                if (i != j) {
                                                                 cplx c1 = (fa[j] + fa[i].conj());
cplx c2 = (fa[j] - fa[i].conj()) * r2;
 constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
 B = (B - A + M2) * r12 % M2;
 C = (C - A + M3) * r13 % M3;
                                                                 cplx d1 = (fb[j] + fb[i].conj()) * r3;
 C = (C - B + M3) * r23 % M3;
                                                                 cplx d2 = (fb[j] - fb[i].conj()) * r4;
                                                                 fa[i] = c1 * d1 + c2 * d2 * r5;
  return (A + B * M1 + C * M1M2) % mod;
                                                                 fb[i] = c1 * d2 + c2 * d1;
                                                                fa[j] = a1 * b1 + a2 * b2 * r5;
namespace fft {
                                                                fb[j] = a1 * b2 + a2 * b1;
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
                                                               fft(fa, sz), fft(fb, sz);
cplx omega[maxn + 1];
                                                               vector<int> res(sz);
void prefft() {
                                                               for (int i = 0; i < sz; ++i) {
for (int i = 0; i <= maxn; i++)</pre>
                                                                long long a = round(fa[i].re), b = round(fb[i].re),
 omega[i] = cplx(cos(2 * pi * j / maxn),
                                                                     c = round(fa[i].im);
     sin(2 * pi * j / maxn));
                                                                res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
void fft(vector<cplx> &v, int n) {
                                                               return res:
int z = __builtin_ctz(n) - 1;
                                                              }}
for (int i = 0; i < n; ++i) {
                                                                   Chinese Remainder
                                                              5 11
 int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^{=(i >> j & 1) << (z - j);
                                                              1ld crt(lld ans[], lld pri[], int n){
 if (x > i) swap(v[x], v[i]);
                                                               lld M = 1, ret = 0;
                                                               for(int i=0;i<n;i++) M *= pri[i];</pre>
for (int s = 2; s <= n; s <<= 1) {
                                                               for(int i=0;i<n;i++){</pre>
                                                                lld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
 int z = s >> 1;
for (int i = 0; i < n; i += s) {</pre>
                                                                ret += (ans[i]*(M/pri[i])%M * iv)%M;
  for (int k = 0; k < z; ++k) {
                                                                ret %= M;
    cplx x = v[i + z + k] * omega[maxn / s * k];
                                                               }
    v[i + z + k] = v[i + k] - x;
                                                               return ret;
    v[i+k] = v[i+k] + x;
                                                              }
                                                              /*
                                                              Another:
                                                              x = a1 \% m1
                                                              x = a2 \% m2
void ifft(vector<cplx> &v, int n) {
                                                              g = gcd(m1, m2)
fft(v, n);
                                                              assert((a1-a2)%g==0)
reverse(v.begin() + 1, v.end());
                                                              [p, q] = exgcd(m2/g, m1/g)
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
                                                              return a2+m2*(p*(a1-a2)/g)
                                                              0 <= x < lcm(m1, m2)
VL convolution(const VI &a, const VI &b) {
// Should be able to handle N <= 10^5, C <= 10^4
                                                              5.12 Berlekamp Massey
int sz = 1;
while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
                                                              // x: 1-base, p[]: 0-base
vector<cplx> v(sz);
                                                              template<size_t N>
for (int i = 0; i < sz; ++i) {
  double re = i < a.size() ? a[i] : 0;</pre>
                                                              vector<llf> BM(llf x[N], size_t n){
                                                               size_t f[N]={0},t=0;11f d[N];
  double im = i < b.size() ? b[i] : 0;</pre>
                                                               vector<llf> p[N];
 v[i] = cplx(re, im);
                                                               for(size_t i=1,b=0;i<=n;++i) {</pre>
                                                                for(size_t j=0;j<p[t].size();++j)</pre>
fft(v, sz);
                                                                 d[i]+=x[i-j-1]*p[t][j];
for (int i = 0; i \le sz / 2; ++i) {
                                                                if(abs(d[i]-=x[i])<=EPS)continue;</pre>
 int j = (sz - i) & (sz - 1);
                                                                f[t]=i;if(!t){p[++t].resize(i);continue;}
 cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
                                                                vector<llf> cur(i-f[b]-1);
  * cplx(0, -0.25);
if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
                                                                11f k=-d[i]/d[f[b]];cur.PB(-k);
                                                                for(size_t j=0;j<p[b].size();j++)</pre>
    ].conj()) * cplx(0, -0.25);
                                                                 cur.PB(p[b][j]*k);
  v[i] = x;
                                                                if(cur.size()<p[t].size())cur.resize(p[t].size());</pre>
                                                                for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];</pre>
ifft(v, sz);
                                                                if(i-f[b]+p[b].size()>=p[t].size()) b=t;
```

```
Poly X(*this, _n), Y(rhs, _n);
ntt(X.data(), _n), ntt(Y.data(), _n);
 p[++t]=cur;
                                                                  fi(0, _n) X[i] = X[i] * Y[i] % P;
return p[t];
                                                                  ntt(X.data(), _n, true);
                                                                  return X.isz(n() + rhs.n() - 1);
5.13 NTT
template <int mod, int G, int maxn>
                                                                 Poly Inv() const { // coef[0] != 0
struct NTT {
                                                                  if (n() == 1) return {ntt.minv(coef[0])};
static_assert (maxn == (maxn & -maxn));
                                                                  const int _n = n2k(n() * 2);
                                                                  Poly Xi = Poly(*this, (n() + 1)/2).Inv().isz(_n);
 int roots[maxn];
                                                                  Poly Y(*this, _n);
NTT () {
 int r = modpow(G, (mod - 1) / maxn);
                                                                  ntt(Xi.data(), _n), ntt(Y.data(), _n);
                                                                  fi(0, _n) {
Xi[i] *= (2 - Xi[i] * Y[i]) % P;
 for (int i = maxn >> 1; i; i >>= 1) {
   roots[i] = 1;
   for (int j = 1; j < i; j++)
                                                                   if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
    roots[i + j] = modmul(roots[i + j - 1], r);
   r = modmul(r, r);
                                                                  ntt(Xi.data(), _n, true);
 }
                                                                  return Xi.isz(n());
 // n must be 2^k, and 0 \le F[i] < mod
                                                                 Poly Sqrt() const { // Jacobi(coef[0], P) = 1
void inplace_ntt(int n, int F[], bool inv = false) {
                                                                  if (n()==1) return {QuadraticResidue(coef[0], P)};
 for (int i = 0, j = 0; i < n; i++) {
  if (i < j) swap(F[i], F[j]);
  for (int k = n>>1; (j^=k) < k; k>>=1);
                                                                  Poly X = Poly(*this, (n()+1) / 2).Sqrt().isz(n());
                                                                  return X.iadd(Mul(X.Inv()).isz(n())).imul(P/2+1);
                                                                 pair<Poly, Poly> DivMod(const Poly &rhs) const {
  for (int s = 1; s < n; s *= 2) {
                                                                  // (rhs.)back() != 0
   for (int i = 0; i < n; i += s * 2) {
                                                                  if (n() < rhs.n()) return {{0}, *this};</pre>
    for (int j = 0; j < s; j++) {
                                                                  const int _n = n() - rhs.n() + 1;
                                                                  Poly X(rhs); X.irev().isz(_n);
Poly Y(*this); Y.irev().isz(_n);
     int a = F[i+j];
     int b = modmul(F[i+j+s], roots[s+j]);
                                                                  Poly Q = Y.Mul(X.Inv()).isz(_n).irev();
     F[i+j] = modadd(a, b); // a + b
     F[i+j+s] = modsub(a, b); // a - b
                                                                  X = rhs.Mul(Q), Y = *this;
fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
                                                                  return {Q, Y.isz(max(1, rhs.n() - 1))};
   }
  if (inv) {
                                                                 Poly Dx() const {
                                                                  Poly ret(n() - \dot{1});
  int invn = modinv(n);
                                                                  fi(0, ret.n()) ret[i] = (i + 1) * coef[i + 1] % P;
   for (int i = 0; i < n; i++)
    F[i] = modmul(F[i], invn);
                                                                  return ret.isz(max(1, ret.n()));
   reverse(F + 1, F + n);
                                                                 Poly Sx() const {
                                                                  Poly ret(n() + 1);
                                                                  fi(0, n()) ret[i + 1]=ntt.minv(i + 1)*coef[i] % P;
const int P=2013265921, root=31;
                                                                  return ret:
const int MAXN=1<<20;</pre>
NTT<P, root, MAXN> ntt;
                                                                 Poly _tmul(int nn, const Poly &rhs) const {
                                                                  Poly Y = Mul(rhs).isz(n() + nn - 1);
      Polynomial Operations
                                                                  return Poly(Y.data() + n() - 1, nn);
using VL = vector<LL>
#define fi(s, n) for (int i=int(s); i<int(n); ++i)</pre>
                                                                 VL _eval(const VL &x, const auto up)const{
                                                                  const int _n = (int)x.size();
#define Fi(s, n) for (int i=int(n); i>int(s); --i)
int n2k(int n) {
                                                                  if (!_n) return {};
                                                                  vector<Poly> down(_n * 2);
int sz = 1; while (sz < n) sz <<= 1;</pre>
                                                                  down[1] = DivMod(up[1]).second;
return sz;
                                                                  fi(2,_n*2) down[i]=down[i/2].DivMod(up[i]).second;
                                                                  /* down[1] = Poly(up[1]).irev().isz(n()).Inv().irev()
template<int MAXN, LL P, LL RT> // MAXN = 2^k
struct Poly { // coefficients in [0, P)
                                                                      _tmul(_n, *this);
                                                                  fi(2, _n * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
    1, down[i / 2]); */
static NTT<MAXN, P, RT> ntt;
VL coef;
                                                                  VL_y(_n);
int n() const { return coef.size(); } // n()>=1
                                                                  fi(0, _n) y[i] = down[_n + i][0];
LL *data() { return coef.data(); }
const LL *data() const { return coef.data(); }
                                                                  return y;
LL &operator[](size_t i) { return coef[i]; }
                                                                 static vector<Poly> _tree1(const VL &x) {
 const LL &operator[](size_t i)const{return coef[i];}
Poly(initializer_list<LL> a) : coef(a) { }
                                                                  const int _n = (int)x.size();
 explicit Poly(int _n = 1) : coef(_n) { }
                                                                  vector<Poly> up(_n * 2);
Poly(const LL *arr, int _n) : coef(arr, arr + _n) {}
Poly(const Poly &p, int _n) : coef(_n) {
    copy_n(p.data(), min(p.n(), _n), data());
                                                                  return up;
Poly& irev(){return reverse(data(),data()+n()),*this;}
Poly& isz(int _n) { return coef.resize(_n), *this; }
                                                                 VL Eval(const VL&x)const{return _eval(x,_tree1(x));}
                                                                 static Poly Interpolate(const VL &x, const VL &y) {
Poly& iadd(const Poly &rhs) { // n() == rhs.n()
                                                                  const int _n = (int)x.size();
                                                                  vector<Poly> up = _tree1(x), down(_n * 2);
VL z = up[1].Dx()._eval(x, up);
 fi(0, n()) if ((coef[i]+=rhs[i]) >= P)coef[i]-=P;
  return *this;
                                                                  fi(0, _n) z[i] = y[i] * ntt.minv(z[i]) % P;
                                                                  fi(0, _n) down[_n + i] = \{z[i]\};

Fi(0, _n-1) down[i] = down[i * 2].Mul(up[i * 2 + 1])
Poly& imul(LL k) {
 fi(0, n()) coef[i] = coef[i] * k % P;
                                                                   .iadd(down[i * 2 + 1].Mul(up[i * 2]));
  return *this;
                                                                  return down[1];
Poly Mul(const Poly &rhs) const {
                                                                 Poly Ln() const { // coef[0] == 1
 const int _n = n2k(n() + rhs.n() - 1);
```

```
return Dx().Mul(Inv()).Sx().isz(n());
                                                                    if( inv )
                                                                     for( int i = 0 ; i < N ; i++ ) {
 Poly Exp() const \{ // coef[0] == 0 \}
  if (n() == 1) return {1};
Poly X = Poly(*this, (n() + 1)/2).Exp().isz(n());
                                                                      x[ i ] *= inv( N, MOD );
                                                                      x[ i ] %= MOD:
  Poly Y = X.Ln(); Y[0] = P - 1;
                                                                     }
  fi(0, n()) if((Y[i] = coef[i] - Y[i]) < 0)Y[i]+=P;
  return X.Mul(Y).isz(n());
                                                                          DiscreteLog
                                                                   5.16
 Poly Pow(const string &K) const {
  int nz = 0;
                                                                   template<typename Int>
  while (nz < n() && !coef[nz]) ++nz;</pre>
                                                                   Int BSGS(Int x, Int y, Int M) {
  LL nk = 0, nk2 = 0;
                                                                     // x^? \equiv y (mod M)
  for (char c : K) {
                                                                     Int t = 1, c = 0, g = 1;
   nk = (nk * 10 + c - '0') % P;
                                                                     for (Int M_ = M; M_ > 0; M_ >>= 1)
   nk2 = nk2 * 10 + c - '0';
                                                                       g = g * x % M;
   if (nk2 * nz >= n()) return Poly(n());
                                                                     for (g = gcd(g, M); t % g != 0; ++c) {
   nk2 %= P - 1:
                                                                       if (t == y) return c;
                                                                        t = t * x % M;
  if (!nk && !nk2) return Poly({1}, n());
  Poly X(data() + nz, n() - nz * nk2);
                                                                     if (y % g != 0) return -1;
                                                                     t /= g, y /= g, M /= g;
  LL x0 = X[0];
                                                                     Int h = 0, gs = 1;
for (; h * h < M; ++h) gs = gs * x % M;
  return X.imul(ntt.minv(x0)).Ln().imul(nk).Exp()
   .imul(ntt.mpow(x0, nk2)).irev().isz(n()).irev();
                                                                     unordered_map<Int, Int> bs;
 Poly InvMod(int L) { // (to evaluate linear recursion)
                                                                     for (Int s = 0; s < h; bs[y] = ++s)
  Poly R{1, \theta}; // *this * R mod x^L = 1 (*this[\theta] ==
                                                                       y = y * x % M;
                                                                     for (Int s = 0; s < M; s += h) {
  for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                                       t = t * gs % M;
   Poly 0 = R.Mul(Poly(data(), min(2 << level, n())));
                                                                       if (bs.count(t)) return c + s + h - bs[t];
   Poly Q(2 << level); Q[0] = 1;
   for (int j = (1 << level); j < (2 << level); ++j)
Q[j] = (P - O[j]) % P;</pre>
                                                                     return -1;
   R = R.Mul(Q).isz(4 << level);
  }
                                                                   5.17
                                                                          FloorSum
  return R.isz(L);
                                                                   // @param n `n < 2^32`
 static LL LinearRecursion(const VL&a,const VL&c,LL n){
                                                                   // @param m `1 <= m < 2^32`
  // a_n = \sum_{j=0}^{n} a_{n-j}
                                                                   // @return sum_\{i=0\}^{n-1} floor((ai + b)/m) mod 2^64
  const int k = (int)a.size();
                                                                   1lu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
  assert((int)c.size() == k + 1);
                                                                    11u ans = 0;
  Poly C(k + 1), W(\{1\}, k), M = \{0, 1\}; fi(1, k + 1) C[k - i] = c[i] ? P - c[i] : 0;
                                                                    while (true)
                                                                     if (a >= m) {
                                                                      ans += n * (n - 1) / 2 * (a / m); a %= m;
  while (n) {
   if (n % 2) W = W.Mul(M).DivMod(C).second;
                                                                     if (b >= m) {
   n /= 2, M = M.Mul(M).DivMod(C).second;
                                                                      ans += n * (b / m); b %= m;
  LL ret = 0;
                                                                     llu y_max = a * n + b;
  fi(0, k) ret = (ret + W[i] * a[i]) % P;
                                                                     if (y_max < m) break;</pre>
  return ret;
                                                                     // y_max < m * (n + 1)
                                                                     // floor(y_max / m) <= n
};
                                                                     n = (1lu)(y_max / m), b = (1lu)(y_max % m);
#undef fi
                                                                     swap(m, a);
#undef Fi
using Poly_t = Poly<131072 * 2, 998244353, 3>:
                                                                    return ans;
template<> decltype(Poly_t::ntt) Poly_t::ntt = {};
                                                                   11d floor_sum(11d n, 11d m, 11d a, 11d b) {
5.15 FWT
                                                                    llu ans = 0;
/* xor convolution:
                                                                    if (a < 0) {
 * x = (x0, x1) , y = (y0, y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                     11u \ a2 = (a \% m + m) \% m;
                                                                     ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
                                                                     a = a2:
* x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)

* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))

* z = (1/2) * z''
                                                                    if (b < 0) {
                                                                     11u b2 = (b \% m + m) \% m;
 * or convolution:
                                                                     ans -= 1ULL * n * ((b2 - b) / m);
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                                     b = b2:
 * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
                                                                    return ans + floor_sum_unsigned(n, m, a, b);
const LL MOD = 1e9+7;
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
 for( int d = 1 ; d < N ; d <<= 1 ) {
                                                                          ExtendedFloorSum
                                                                   5.18
  int d2 = d << 1;
  for( int s = 0 ; s < N ; s += d2 )
   for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
  LL ta = x[ i ] , tb = x[ j ];
  x[ i ] = ta+tb;</pre>
                                                                    g(a, b, c, n) = \sum_{i=1}^{n} i \lfloor \frac{ai + b}{c} \rfloor
                                                                                 \left( \left| \frac{a}{c} \right| \cdot \frac{n(n+1)(2n+1)}{6} + \left| \frac{b}{c} \right| \cdot \frac{n(n+1)}{2} \right)
    x[ j ] = ta-tb;
                                                                                  +g(a \bmod c, b \bmod c, c, n),
                                                                                                                      a \geq c \vee b \geq c
    if('x[ i ] >= MOD ) x[ i ] -= MOD;
                                                                                                                      n < 0 \lor a = 0
    if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
                                                                                  \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
                                                                                 -h(c, c-b-1, a, m-1)),
                                                                                                                      otherwise
```

```
\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}
```

5.19 Quadratic residue

```
struct S {
 int MOD, w;
 int64_t x, y;
 S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
   : MOD(m), w(w_{-}), x(x_{-}), y(y_{-}) {}
 S operator*(const S &rhs) const {
  int w_ = w;
  if (w_ == -1) w_ = rhs.w;
  assert(w_! = -1 \text{ and } w_ == rhs.w);
  return { MOD, w_,
(x * rhs.x + y * rhs.y % MOD * w) % MOD,
   (x * rhs.y + y * rhs.x) % MOD };
int get_root(int n, int P) {
  if (P == 2 or n == 0) return n;
  if (qpow(n, (P - 1) / 2, P) != 1) return -1;
  auto check = [&](int x) {
  return qpow(x, (P - 1) / 2, P); };
  if (check(n) == P-1) return -1
  int64_t a; int w; mt19937 rnd(7122);
  do { a = rnd() % P;
  w = ((a * a - n) % P + P) % P;
  } while (check(w) != P - 1);
  return qpow(S(P, w, a, 1), (P + 1) / 2).x;
```

5.20 De-Bruijn

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
 if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
 } else {
 aux[t] = aux[t - p];
 db(t + 1, p, n, k);
for (int i = aux[t - p] + 1; i < k; ++i) {
  aux[t] = i;
   db(t + 1, t, n, k);
int de_bruijn(int k, int n) {
// return cyclic string of len k^n s.t. every string
// of len n using k char appears as a substring.
if (k == 1) {
 res[0] = 0;
  return 1;
for (int i = 0; i < k * n; i++) aux[i] = 0;
sz = 0;
db(1, 1, n, k);
return sz;
```

5.21 Simplex Construction

Standard form: maximize $\sum_{1\leq i\leq n}c_ix_i$ such that for all $1\leq j\leq m$, $\sum_{1\leq i\leq n}A_{ji}x_i\leq b_j$ and $x_i\geq 0$ for all $1\leq i\leq n$.

- 1. In case of minimization, let $c'_i = -c_i$
- 2. $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j$
- 3. $\sum_{1 \leq i \leq n} A_{ji} x_i = b_j$
 - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$
 - $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

```
5.22 Simplex
```

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return VD(n, -inf) if the solution doesn't exist // return VD(n, +inf) if the solution is unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
 double inv = 1.0 / d[r][s];
 for (int i = 0; i < m + 2; ++i)
  for (int j = 0; j < n + 2; ++j)
if (i != r && j != s)
d[i][j] -= d[r][j] * d[i][s] * inv;</pre>
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
 d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
 int x = m + z;
 while (true) {
  int s = -1;
  for (int i = 0; i <= n; ++i) {
   if (!z && q[i] == -1) continue
   if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
  if (d[x][s] > -eps) return true;
  for (int i = 0; i < m; ++i) {
  if (d[i][s] < eps) continue;</pre>
     d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
  pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
m = b.size(), n = c.size();
d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)</pre>
 p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
 q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
 for (int i = 1; i < m; ++i)
  if (d[i][n + 1] < d[r][n + 1]) r = i;
 if (d[r][n + 1] < -eps) {</pre>
  pivot(r, n);
  if (!phase(1) || d[m + 1][n + 1] < -eps)</pre>
   return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
   int s = min_element(d[i].begin(), d[i].end() - 1)
         - d[i].begin();
   pivot(i, s);
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)</pre>
  if (p[i] < n) x[p[i]] = d[i][n + 1];
 return x;
5.23 Charateristic Polynomial
```

```
vector<vector<int>>> &A) {
  int N = A.size();
  vector<vector<int>>> h = A;
  for (int i = 0; i < N - 2; ++i) {
    if (!H[i + 1][i]) {
      for (int j = i + 2; j < N; ++j) {
        if (H[j][i]) {
            for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j][k]);
        }
}</pre>
```

```
for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k
    ][j]);
     break:
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
     * H[i + 1][k] * (kP - coef)) % kP;
   for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
 return H;
vector<int> CharacteristicPoly(const vector<vector<int</pre>
    >> &A) {
 int N = A.size();
 auto H = Hessenberg(A);
 for (int i = 0; i < N; ++i) {
  for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
 vector<vector<int>>> P(N + 1, vector<int>(N + 1));
 P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
  P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
    1];
  int val = 1;
  for (int j = i - 1; j >= 0; --j) {
   int coef = 1LL * val * H[j][i - 1] % kP;
   for (int k = 0; k \le j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
  }
 if (N & 1) {
  for (int i = 0; i \le N; ++i) P[N][i] = kP - P[N][i];
 return P[N];
5.24
        Partition Number
int b = sqrt(n);
ans[0] = tmp[0] = 1;
for (int i = 1; i <= b; i++) {
 for (int rep = 0; rep < 2; rep++)</pre>
  for (int j = i; j <= n - i * i; j++)
modadd(tmp[j], tmp[j-i]);</pre>
 for (int j = i * i; j <= n; j++)
  modadd(ans[j], tmp[j - i * i]);
6
     Geometry
     Basic Geometru
using coord_t = int;
using Real = double;
using Point = std::complex<coord_t>;
int sgn(coord_t x) {
return (x > 0) - (x < 0);
coord_t dot(Point a, Point b) {
return real(conj(a) * b);
```

coord_t cross(Point a, Point b) {
 return imag(conj(a) * b);

int ori(Point a, Point b, Point c) {

bool operator<(const Point &a, const Point &b) {</pre>

? real(a) < real(b) : imag(a) < imag(b);</pre>

return sgn(cross(b - a, c - a));

return real(a) != real(b)

int qa = (imag(a) == 0

int argCmp(Point a, Point b) {
 // -1 / 0 / 1 <-> < / == / > (atan2)

```
? (real(a) < 0 ? 3 : 1) : (imag(a) < 0 ? 0 : 2));
 int qb = (imag(b) == 0
   ? (real(b) < 0 ? 3 : 1) : (imag(b) < 0 ? 0 : 2));
 if (qa != qb)
  return sgn(qa - qb);
 return sgn(cross(b, a));
template <typename V> Real area(const V & pt) {
 coord_t ret = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
 ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 return ret / 2.0;
6.2 2D Convex Hull
template<typename PT>
vector<PT> buildConvexHull(vector<PT> d) {
 sort(ALL(d), [](const PT& a, const PT& b){
   return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
 vector<PT> s(SZ(d)<<1);</pre>
 int o = 0;
 for(auto p: d) {
  while(o \ge 2 && cross(p - s[o - 2], s[o - 1] - s[o - 2]) <= 0)
  s[o++] = p;
 for(int i=SZ(d)-2, t = o+1; i>=0; i--){
  while(o>=t&&cross(d[i]-s[o-2],s[o-1]-s[o-2])<=0)</pre>
  s[o++] = d[i];
 s.resize(o-1);
 return s:
6.3
     3D Convex Hull
// return the faces with pt indexes
int flag[MXN][MXN];
struct Point{
 ld x,y,z;
 Point operator * (const ld &b) const {
  return (Point) {x*b, y*b, z*b};}
 Point operator * (const Point &b) const {
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
Point ver(Point a, Point b, Point c) {
 return (b - a) * (c - a);}
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0;
 REP(i,n) REP(j,n) flag[i][j] = 0;
 vector<Face> now;
 now.emplace_back(0,1,2);
 now.emplace_back(2,1,0)
 for (int i=3; i<n; i++){
  ftop++; vector<Face> next;
  REP(j, SZ(now)) {
  Face& f=now[j]; int ff = 0;
   ld d=(pt[i]-pt[f.a]).dot(
     ver(pt[f.a], pt[f.b], pt[f.c]));
   if (d <= 0) next.push_back(f);</pre>
   if (d > 0) ff=ftop;
   else if (d < 0) ff=-ftop
   flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
  REP(j, SZ(now)) {
  Face& f=now[j];
   if (flag[f.a][f.b] > 0 &&
     flag[f.a][f.b] != flag[f.b][f.a])
    next.emplace_back(f.a,f.b,i);
   if (flag[f.b][f.c] > 0 &&
     flag[f.b][f.c] \stackrel{!}{=} flag[f.c][f.b])
    next.emplace_back(f.b,f.c,i);
   if (flag[f.c][f.a] > 0 &&
     flag[f.c][f.a] != flag[f.a][f.c])
    next.emplace_back(f.c,f.a,i);
  now=next;
```

return now;

```
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      2D Farthest Pair
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {
  while(abs(cross(stk[i+1]-stk[i],</pre>
   stk[(pos+1)%n]-stk[i])) >
   abs(cross(stk[i+1]-stk[i],
stk[pos]-stk[i]))) pos = (pos+1)%n;
ans = max({ans, dis(stk[i], stk[pos]),
  dis(stk[i+1], stk[pos])});
6.5 2D Closest Pair
struct cmp_y {
bool operator()(const P& p, const P& q) const {
  return p.y < q.y;</pre>
}:
multiset<P, cmp_y> s;
void solve(P a[], int n) {
 sort(a, a + n, [](const P& p, const P& q) {
  return tie(p.x, p.y) < tie(q.x, q.y);</pre>
 llf d = INF; int pt = 0;
 for (int i = 0; i < n; ++i) {
 while (pt < i \text{ and } a[i].x - a[pt].x >= d)
   s.erase(s.find(a[pt++]));
  auto it = s.lower_bound(P(a[i].x, a[i].y - d));
  while (it != s.end() and it->y - a[i].y < d)
   d = min(d, dis(*(it++), a[i]));
  s.insert(a[i]);
 }
}
6.6 kD Closest Pair (3D ver.)
11f solve(vector<P> v) {
 shuffle(v.begin(), v.end(), mt19937());
 unordered_map<lld, unordered_map<lld,
  unordered_map<lld, int>>> m;
 llf d = dis(v[0], v[1]);
 auto Idx = [&d] (llf x) -> lld {
 return round(x * 2 / d) + 0.1; };
auto rebuild_m = [&m, &v, &Idx](int k) {
  m.clear();
  for (int i = 0; i < k; ++i)
m[Idx(v[i].x)][Idx(v[i].y)]</pre>
    [Idx(v[i].z)] = i;
 }; rebuild_m(2);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
     kz = Idx(v[i].z); bool found = false;
  for (int dx = -2; dx <= 2; ++dx) {
   const 11d nx = dx + kx
   if (m.find(nx) == m.end()) continue;
   auto& mm = m[nx];
for (int dy = -2; dy <= 2; ++dy) {</pre>
    const 11d ny = dy + ky;
    if (mm.find(ny) == mm.end()) continue;
    auto& mmm = mm[ny];
    for (int dz = -2; dz <= 2; ++dz) {
     const 1ld nz = dz + kz;
     if (mmm.find(nz) == mmm.end()) continue;
     const int p = mmm[nz];
if (dis(v[p], v[i]) < d) {
  d = dis(v[p], v[i]);</pre>
      found = true;
     }
  if (found) rebuild_m(i + 1);
  else m[kx][ky][kz] = i;
 return d;
6.7 Simulated Annealing
11f anneal() {
 mt19937 rnd_engine( seed );
uniform_real_distribution< llf > rnd(0, 1);
```

```
const 11f dT = 0.001;
 // Argument p
 11f S_cur = calc( p ), S_best = S_cur;
 for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
  // Modify p to p_prime
  const llf S_prime = calc( p_prime );
  const llf delta_c = S_prime - S_cur;
llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
  if ( rnd( rnd_engine ) <= prob )</pre>
  S_cur = S_prime, p = p_prime;
if ( S_prime < S_best ) // find min
   S_best = S_prime, p_best = p_prime;
 return S_best;
}
6.8 Half Plane Intersection
// NOTE: Point is complex<Real>
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
struct Line {
  Point st, ed;
  Point dir
  Line (Point _s, Point _e)
   : st(_s), ed(_e), dir(_e - _s) {}
bool operator<(const Line &lhs, const Line &rhs) {</pre>
  if (int cmp = argCmp(lhs.dir, rhs.dir))
    return cmp == -1;
  return ori(lhs.st, lhs.ed, rhs.st) < 0;</pre>
Point intersect(const Line &A, const Line &B) {
  Real t = cross(B.st - A.st, B.dir) /
   cross(A.dir, B.dir);
  return A.st + t * A.dir;
}
Real HPI(vector<Line> &lines) {
  sort(lines.begin(), lines.end());
  deque<Line> que;
  deque<Point> pt;
  que.push_back(lines[0]);
  for (int i = 1; i < (int)lines.size(); i++) {</pre>
    if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
     continue
#define POP(L, R) \
    while (pt.size() > 0 \
      && ori(L.st, L.ed, pt.back()) < 0) \
      pt.pop_back(), que.pop_back(); \
    while (pt.size() > 0 \
      && ori(R.st, R.ed, pt.front()) < 0) \
      pt.pop_front(), que.pop_front();
    POP(lines[i], lines[i])
    pt.push_back(intersect(que.back(), lines[i]));
    que.push_back(lines[i]);
  POP(que.front(), que.back())
  if (que.size() <= 1 ||</pre>
    argCmp(que.front().dir, que.back().dir) == 0)
  pt.push_back(intersect(que.front(), que.back()));
  return area(pt);
     Minkowski Sum
vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
 hull(A), hull(B);
 vector<pll> C(1, A[0] + B[0]), s1, s2;
for(int i = 0; i < SZ(A); ++i)</pre>
  s1.pb(A[(i + 1) % SZ(A)] - A[i]);
 for(int i = 0; i < SZ(B); i++)
s2.pb(B[(i + 1) % SZ(B)] - B[i]);
 for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
  if (p2 >= SZ(B)
    | | (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
   C.pb(C.back() + s1[p1++]);
  else
   C.pb(C.back() + s2[p2++]);
 return hull(C), C;
```

6.10 Intersection of line and Circle

6.11 Intersection of Polygon and Circle

```
// Divides into multiple triangle, and sum up
// test by HDU2892
const double PI=acos(-1);
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
if(abs(pb)<eps) return 0;</pre>
double S, h, theta;
double a=abs(pb), b=abs(pa), c=abs(pb-pa);
double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
double cosC = dot(pa,pb) / a / b, C = acos(cosC);
if(a > r){
 S = (C/2)*r*r;
 h = a*b*sin(C)/c;
  if (h < r && B < PI/2)</pre>
   S = (acos(h/r)*r*r - h*sqrt(r*r-h*h));
else if(b > r){
 theta = PI - B - asin(sin(B)/r*a);
  S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
else S = .5*sin(C)*a*b;
return S;
double area_poly_circle(const vector<pdd> poly,
 const pdd &0,const double r){
double S=0;
 for(int i=0;i<SZ(poly);++i)</pre>
  S+=_area(poly[i]-0,poly[(i+1)%SZ(poly)]-0,r)
    *ori(0,poly[i],poly[(i+1)%SZ(poly)]);
return fabs(S);
```

6.12 Intersection of Two Circle

6.13 Tangent line of Two Circle

```
vector<Line> go(const Cir& c1,
  const Cir& c2, int sign1){
 // sign1 = 1 for outer tang, -1 for inter tang
vector<Line> ret;
double d_sq = norm2( c1.0 - c2.0 );
if( d_sq < eps ) return ret;</pre>
double d = sqrt( d_sq );
Pt v = (c2.0 - c1.0) / d;
double c = ( c1.R - sign1 * c2.R ) / d;
 if( c * c > 1 ) return ret;
double h = sqrt( max( 0.0 , 1.0 - c * c ) );
for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
 Pt n = \{ v.X * c - sign2 * h * v.Y ,
   v.Y * c + sign2 * h * v.X };
 Pt p1 = c1.0 + n * c1.R;
 Pt p2 = c2.0 + n * (c2.R * sign1);
  if( fabs( p1.X - p2.X ) < eps and
    fabs( p1.Y - p2.Y ) < eps )
   p2 = p1 + perp(c2.0 - c1.0);
```

```
ret.push_back( { p1 , p2 } );
 return ret;
        Minimum Covering Circle
template<typename P>
Circle getCircum(const P &a, const P &b, const P &c){
 Real a1 = a.x-b.x, b1 = a.y-b.y;
 Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
 Real a2 = a.x-c.x, b2 = a.y-c.y;
 Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
 Circle cc;
 cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
 cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
 cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
 return cc:
template<typename P>
Circle MinCircleCover(const vector<P>& pts){
 random_shuffle(pts.begin(), pts.end());
 Circle c = { pts[0], 0 }
 for(int i=0;i<(int)pts.size();i++){</pre>
  if (dist(pts[i], c.o) <= c.r) continue;</pre>
  c = { pts[i], 0 };
for (int j = 0; j < i; j++) {
   if(dist(pts[j], c.o) <= c.r) continue;</pre>
   c.o = (pts[i] + pts[j]) / 2;
   c.r = dist(pts[i], c.o);
   for (int k = 0; k < j; k++) {
    if (dist(pts[k], c.o) <= c.r) continue;</pre>
    c = getCircum(pts[i], pts[j], pts[k]);
  }
 return c;
6.15 KDTree (Nearest Point)
const int MXN = 100005;
struct KDTree {
 struct Node {
  int x,y,x1,y1,x2,y2;
  int id,f;
Node *L, *R;
 } tree[MXN], *root;
 int n:
 LL dis2(int x1, int y1, int x2, int y2) {
  LL dx = x1-x2, dy = y1-y2;
  return dx*dx+dy*dy;
 static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
 static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 void init(vector<pair<int,int>> ip) {
  n = ip.size();
  for (int i=0; i<n; i++) {</pre>
   tree[i].id = i;
   tree[i].x = ip[i].first;
   tree[i].y = ip[i].second;
  root = build_tree(0, n-1, 0);
 Node* build_tree(int L, int R, int d) {
  if (L>R) return nullptr;
  int M = (L+R)/2; tree[M].f = d%2;
  nth_element(tree+L,tree+M,tree+R+1,d%2?cmpy:cmpx);
  tree[M].x1 = tree[M].x2 = tree[M].x;
  tree[M].y1 = tree[M].y2 = tree[M].y;
  tree[M].L = build_tree(L, M-1, d+1);
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
   tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
  tree[M].R = build_tree(M+1, R, d+1);
  if (tree[M].R) {
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
   tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
```

bool disjuct(Cir &a, Cir &b, int x)

```
{return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
                                                               bool contain(Cir &a, Cir &b, int x)
  return tree+M;
                                                               {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
                                                               bool contain(int i, int j) {
  /* c[j] is non-strictly in c[i]. */
 int touch(Node* r, int x, int y, LL d2){
  LL dis = sqrt(d2)+1;
  if (x<r->x1-dis || x>r->x2+dis ||
                                                                return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c[i].R) \mid c[i].R - c[i].R)
                                                                   [j].R) == 0 && i < j)) && contain(c[i], c[j], -1);
    y<r->y1-dis || y>r->y2+dis)
   return 0:
                                                               void solve(){
  return 1;
                                                                fill_n(Area, C + 2, 0);
                                                                for(int i = 0; i < C; ++i)
 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
  if (!r || !touch(r, x, y, md2)) return;
LL d2 = dis2(r->x, r->y, x, y);
                                                                 for(int j = 0; j < C; ++j)
                                                                  overlap[i][j] = contain(i, j);
  if (d2 < md2 \mid \mid (d2 == md2 \&\& mID < r->id)) {
                                                                for(int i = 0; i < C; ++i)</pre>
   mID = r -> id;
                                                                 for(int j = 0; j < C; ++j)
                                                                  g[i][j] = !(overlap[i][j] || overlap[j][i] ||
   md2 = d2:
                                                                     disjuct(c[i], c[j], -1));
                                                                for(int i = 0; i < C; ++i){</pre>
  // search order depends on split dim
  if ((r->f == 0 && x < r->x) ||
                                                                 int E = 0, cnt = 1;
    (r->f == 1 \&\& y < r->y)) {
                                                                 for(int j = 0; j < C; ++j)
                                                                  if(j != i && overlap[j][i])
   nearest(r->L, x, y, mID, md2);
                                                                   ++cnt;
   nearest(r->R, x, y, mID, md2);
                                                                 for(int j = 0; j < C; ++j)
  } else {
                                                                  if(i != j && g[i][j]) {
   nearest(r->R, x, y, mID, md2);
   nearest(r->L, x, y, mID, md2);
                                                                   pdd aa, bb;
                                                                   CCinter(c[i], c[j], aa, bb);
double A = atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X)
  }
 }
 int query(int x, int y) {
  int id = 1029384756;
                                                                   double B = atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X)
  LL d2 = 102938475612345678LL;
  nearest(root, x, y, id, d2);
                                                                   eve[E++] = Teve(bb, B, 1), eve[E++] = Teve(aa, A,
  return id;
                                                                   -1)
                                                                   if(B > A) ++cnt;
} tree;
                                                                 if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
      Rotating Sweep Line
                                                                 else{
                                                                  sort(eve, eve + E);
void rotatingSweepLine(pair<int, int> a[], int n) {
                                                                  eve[E] = eve[0];
 vector<pair<int, int>> 1;
 1.reserve(n * (n - 1) / 2);
                                                                  for(int j = 0; j < E; ++j){
 for (int i = 0; i < n; ++i)
                                                                   cnt += eve[j].add;
  for (int j = i + 1; j < n; ++j)
                                                                   Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
   1.emplace_back(i, j);
                                                                   double theta = eve[j + 1].ang - eve[j].ang;
                                                                   if (theta < 0) theta += 2. * pi;</pre>
 sort(1.begin(), 1.end(), [&a](auto &u, auto &v){
  1ld udx = a[u.first].first - a[u.second].first;
                                                                   Area[cnt] += (theta - sin(theta)) * c[i].R * c[i].
  11d udy = a[u.first].second - a[u.second].second;
                                                                  R * .5;
  1ld vdx = a[v.first].first - a[v.second].first;
                                                                  }
  11d vdy = a[v.first].second - a[v.second].second;
  if (udx == 0 \text{ or } vdx == 0) \text{ return not } udx == 0;
  int s = sgn(udx * vdx);
                                                              };
  return udy * vdx * s < vdy * udx * s;
 });
 vector<int> idx(n), p(n);
                                                                   Stringology
 iota(idx.begin(), idx.end(), 0);
 sort(idx.begin(), idx.end(), [&a](int i, int j){
                                                              7.1 Hash
 return a[i] < a[j]; });
for (int i = 0; i < n; ++i) p[idx[i]] = i;
                                                              class Hash {
                                                               private:
 for (auto [i, j]: 1) {
                                                                static constexpr int P = 127, Q = 1051762951;
  // do here
                                                                vector<int> h, p;
  swap(p[i], p[j]);
                                                               public
  idx[p[i]] = i, idx[p[j]] = j;
                                                                void init(const string &s){
                                                                 h.assign(s.size()+1, 0); p.resize(s.size()+1);
                                                                 for (size_t i = 0; i < s.size(); ++i)</pre>
      Circle Cover
                                                                  h[i + 1] = add(mul(h[i], P), s[i]);
6.17
                                                                 generate(p.begin(), p.end(),[x=1,y=1,this]()
const int N = 1021;
                                                                   mutable{y=x;x=mul(x,P);return y;});
struct CircleCover {
 int C;
                                                                int query(int 1, int r){ // 1-base (1, r]}
 Cir c[N]
                                                                 return sub(h[r], mul(h[1], p[r-1]));}
 bool g[N][N], overlap[N][N];
 // Area[i] : area covered by at least i circles
 double Area[ N ];
                                                              7.2 Suffix Array
 void init(int _C){ C = _C;}
 struct Teve {
                                                              namespace sfxarray {
 pdd p; double ang; int add;
Teve() {}
                                                              bool t[maxn * 2];
                                                              int hi[maxn], rev[maxn];
                                                              int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
  Teve(pdd _a, double _b, int _c):p(_a), ang(_b), add(
                                                              int x[maxn], p[maxn], q[maxn * 2];
    _c){}
  bool operator<(const Teve &a)const
                                                              // sa[i]: sa[i]-th suffix is the \
  {return ang < a.ang;}
                                                              // i-th lexigraphically smallest suffix.
 }eve[N * 2];
                                                              // hi[i]: longest common prefix \
 // strict: x = 0, otherwise x = -1
                                                              // of suffix sa[i] and suffix sa[i - 1].
                                                              void pre(int *sa, int *c, int n, int z) {
```

LAST->green = potential_green;

```
memset(sa, 0, sizeof(int) * n);
                                                                    else {
 memcpy(x, c, sizeof(int) * z);
                                                                 //assert(potential_green->max_len>(cursor->max_len+1));
                                                                     Node *wish = new Node((cursor->max_len) + 1);
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
                                                                     for(;cursor && cursor->edge[c]==potential_green;
memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
if (sa[i] && !t[sa[i] - 1])
                                                                        cursor = cursor->green)
                                                                      cursor->edge[c] = wish;
                                                                     for (int i = 0; i < 26; i++)
   sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
                                                                     wish->edge[i] = potential_green->edge[i];
memcpy(x, c, sizeof(int) * z);
for (int i = n - 1; i >= 0; --i)
if (sa[i] && t[sa[i] - 1])
                                                                     wish->green = potential_green->green;
                                                                     potential_green->green = wish;
                                                                     LAST->green = wish;
   sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
                                                                  }
void sais(int *s, int *sa, int *p, int *q,
bool *t, int *c, int n, int z) {
                                                                 char S[10000001], A[10000001];
                                                                 int N;
 bool uniq = t[n - 1] = true;
                                                                 int main(){
 int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
                                                                   scanf("%d%s", &N, S);
 memset(c, 0, sizeof(int) * z);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
                                                                   ROOT = LAST = new Node(0);
                                                                   for (int i = 0; S[i]; i++)
Extend(S[i] - 'a');
 for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
 if (uniq) {
  for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
                                                                   while (N--){
  scanf("%s", A);
                                                                    Node *cursor = ROOT;
                                                                    bool ans = true;
 for (int i = n - 2; i \ge 0; --i)
  t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
                                                                    for (int i = 0; A[i]; i++){
                                                                    cursor = cursor->edge[A[i] - 'a'];
 pre(sa, c, n, z);
 for (int i = 1; i <= n - 1; ++i)</pre>
                                                                     if (!cursor) {
  if (t[i] && !t[i - 1])
                                                                      ans = false;
   sa[--x[s[i]]] = p[q[i] = nn++] = i;
                                                                      break;
induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) {
  if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {</pre>
                                                                   puts(ans ? "Yes" : "No");
  bool neq = last < 0 || \</pre>
   memcmp(s + sa[i], s + last,
(p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
                                                                   return 0;
  ns[q[last = sa[i]]] = nmxz += neq;
                                                                 7.4
                                                                      KMP
 }}
 sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
                                                                 vector<int> kmp(const string &s) {
 pre(sa, c, n, z);
                                                                   vector<int> f(s.size(), 0);
 for (int i = nn - 1; i >= 0; --i)
                                                                   /* f[i] = length of the longest prefix
  sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
                                                                     (excluding s[0:i]) such that it coincides
 induce(sa, c, s, t, n, z);
                                                                     with the suffix of s[0:i] of the same length */
                                                                   /* i + 1 - f[i] is the length of the
void build(const string &s) {
  for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
                                                                     smallest recurring period of s[0:i] */
                                                                   int k = 0;
 _s[(int)s.size()] = 0; // s shouldn't contain 0
                                                                   for (int i = 1; i < (int)s.size(); ++i) {</pre>
 sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];</pre>
                                                                   while (k > 0 \&\& s[i] != s[k]) k = f[k-1];
                                                                    if (s[i] == s[k]) ++k;
 for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
                                                                   f[i] = k;
 int ind = 0; hi[0] = 0;
 for (int i = 0; i < (int)s.size(); ++i) {</pre>
                                                                   return f;
  if (!rev[i]) {
   ind = 0;
                                                                 vector<int> search(const string &s, const string &t) {
   continue;
                                                                   // return 0-indexed occurrence of t in s
                                                                   vector<int> f = kmp(t), r;
for (int i = 0, k = 0; i < (int)s.size(); ++i)</pre>
  while (i + ind < (int)s.size() && \</pre>
   s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
                                                                   while(k > 0 && (k==(int)t.size() \mid \mid s[i]!=t[k]))
  hi[rev[i]] = ind ? ind-- : 0;
                                                                    k = f[k - 1];
                                                                    if (s[i] == t[k]) ++k;
                                                                   if (k == (int)t.size()) r.push_back(i-t.size()+1);
7.3
     Suffix Automaton
                                                                   return res;
struct Node{
 Node *green, *edge[26];
 int max_len;
                                                                       Z value
                                                                 7.5
 Node(const int _max_len)
                                                                 char s[MAXN];
  : green(NULL), max_len(_max_len){
                                                                 int len,z[MAXN];
  memset(edge,0,sizeof(edge));
                                                                 void Z_value() {
                                                                   int i,j,left,right;
} *ROOT, *LAST;
                                                                   z[left=right=0]=len;
void Extend(const int c) {
                                                                   for(i=1;i<len;i++)</pre>
 Node *cursor = LAST;
                                                                    j=max(min(z[i-left],right-i),0);
 LAST = new Node((LAST->max_len) + 1);
                                                                    for(;i+j<len&&s[i+j]==s[j];j++);</pre>
 for(;cursor&&!cursor->edge[c]; cursor=cursor->green)
                                                                    if(i+(z[i]=j)>right)right=i+z[left=i];
  cursor->edge[c] = LAST;
 if (!cursor)
                                                                 }
 LAST->green = ROOT;
 else {
                                                                 7.6
                                                                      Manacher
  Node *potential_green = cursor->edge[c];
  if((potential_green->max_len)==(cursor->max_len+1))
                                                                 int z[maxn];
```

int manacher(const string& s) {

```
string t = ".";
for(char c: s) t += c, t += '.';
int l = 0, r = 0, ans = 0;
for (int i = 1; i < t.length(); ++i) {
  z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
  while (i - z[i] >= 0 && i + z[i] < t.length()) {
    if(t[i - z[i]] == t[i + z[i]]) ++z[i];
    else break;
}
if (i + z[i] > r) r = i + z[i], l = i;
}
for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
return ans;
}</pre>
```

7.7 Lexico Smallest Rotation

```
string mcp(string s){
  int n = s.length();
  s += s;
  int i=0, j=1;
  while (i<n && j<n){
    int k = 0;
    while (k < n && s[i+k] == s[j+k]) k++;
    if (s[i+k] <= s[j+k]) j += k+1;
    else i += k+1;
    if (i == j) j++;
  }
  int ans = i < n ? i : j;
  return s.substr(ans, n);
}</pre>
```

7.8 BWT

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a
vector<int> v[ SIGMA ];
void BWT(char* ori, char* res){
  // make ori -> ori + ori
  // then build suffix array
void iBWT(char* ori, char* res){
  for( int i = 0 ; i < SIGMA ; i ++ )</pre>
   v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )</pre>
   v[ ori[i] - BASE ].push_back( i );
  vector<int> a:
  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
  for( auto j : v[ i ] ){
   a.push_back( j );
}
    ori[ ptr ++ ] = BASE + i;
  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
   ptr = a[ ptr ];
  res[ len ] = 0;
}
} bwt;
```

7.9 Palindromic Tree

```
struct palindromic_tree{
struct node{
 int next[26],f,len;
 int cnt, num, st, ed;
 node(int 1=0):f(0),len(1),cnt(0),num(0) {
  memset(next, 0, sizeof(next)); }
vector<node> st;
vector<char> s;
int last, n;
void init(){
 st.clear();s.clear();last=1; n=0;
 st.push_back(0);st.push_back(-1);
 st[0].f=1;s.push_back(-1); }
int getFail(int x){
 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
 return x;}
void add(int c){
 s.push_back(c-='a'); ++n;
 int cur=getFail(last);
```

```
if(!st[cur].next[c]){
   int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
   st[now].num=st[st[now].f].num+1;
  last=st[cur].next[c];
  ++st[last].cnt;}
 void dpcnt() {
  for (int i=st.size()-1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
 int size(){ return st.size()-2;}
} pt;
int main() {
 string s; cin >> s; pt.init();
 for (int i=0; i<SZ(s); i++) {</pre>
 int prvsz = pt.size(); pt.add(s[i]);
 if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [l,r]: s.substr(l, r-l+1)
 return 0;
```

8 Misc

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\ldots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \le k \le n$.

8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W \subseteq X, |W| \leq |N_G(W)|$

8.1.7 Euler's planar graph formula

```
V - E + F = C + 1, E \le 3V - 6(?)
```

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Lucas's theorem

```
{m\choose n}\equiv\prod_{i=0}^k{m_i\choose n_i}\ (\text{mod }p)\text{, where }m=m_kp^k+m_{k-1}p^{k-1}+\cdots+m_1p+m_0, and n=n_kp^k+n_{k-1}p^{k-1}+\cdots+n_1p+n_0.
```

8.1.10 Matroid Intersection

```
Given matroids M_1=(G,I_1),M_2=(G,I_2), find maximum S\in I_1\cap I_2. For each iteration, build the directed graph and find a shortest path from s to t.
```

```
• s \to x: S \sqcup \{x\} \in I_1

• x \to t: S \sqcup \{x\} \in I_2

• y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1 (y is in the unique circuit of S \sqcup \{x\})

• x \to y: S \setminus \{y\} \sqcup \{x\} \in I_2 (y is in the unique circuit of S \sqcup \{x\})
```

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|.$ In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

8.2 DP-opt Condition

8.2.1 totally monotone (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

8.2.2 monge condition (concave/convex)

```
8.3 Convex 1D/1D DP struct segment {
```

```
int i, 1, r;
segment() {}
segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
dp[0] = 0;
deque<segment> dq; dq.push_back(segment(0, 1, n));
for (int i = 1; i <= n; ++i) {
 dp[i] = f(dq.front().i, i)
 while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
 dq.front().l = i + 1;
 segment seg = segment(i, i + 1, n);
 while (dq.size() &&
  f(i, dq.back().1)<f(dq.back().i, dq.back().1))
    dq.pop_back();
  if (dq.size())
  int d = 1 << 20, c = dq.back().1;</pre>
  while (d >>= 1) if (c + d <= dq.back().r)
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
  dq.back().r = c; seg.l = c + 1;
  if (seg.1 <= n) dq.push_back(seg);</pre>
```

8.4 ConvexHull Optimization

```
struct Line {
 mutable int64_t a, b, p;
 bool operator<(const Line &rhs) const { return a < rhs
 bool operator<(int64_t x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>> {
 static const int64_t kInf = 1e18;
 bool Isect(iterator x, iterator y)
  auto Div = [](int64_t a, int64_t b) {
  return a / b - ((a ^ b) < 0 && a % b); }
if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(int64_t a, int64_t b) {
 auto z = insert({a, b, 0}), y = z++, x = y;
while (Isect(y, z)) z = erase(z);
  if (x != begin() \&\& Isect(--x, y)) Isect(x, y = erase)
    (y))
  while ((y = x) != begin() && (--x)->p >= y->p) Isect(
    x, erase(y));
 int64_t Query(int64_t x) {
 auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
};
```

8.5 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

```
8.6 Cactus Matching
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
 dfn[u]=low[u]=++dfs_idx;
 for(int i=0;i<(int)init_g[u].size();i++){</pre>
  int v=init_g[u][i];
  if(v==par[u]) continue;
  if(!dfn[v]){
  par[v]=u;
   tarjan(v)
   low[u]=min(low[u],low[v]);
   if(dfn[u]<low[v]){</pre>
    g[u].push_back(v)
    g[v].push_back(u);
   }
  }else{
   low[u]=min(low[u],dfn[v]);
   if(dfn[v]<dfn[u]){</pre>
    int temp_v=u;
    bcc_id++
    while(temp_v!=v){
     g[bcc_id+n].push_back(temp_v);
     g[temp_v].push_back(bcc_id+n);
     temp_v=par[temp_v];
    g[bcc_id+n].push_back(v);
    g[v].push_back(bcc_id+n);
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
void dfs(int u,int fa){
if(u<=n){
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
   dfs(v,u);
   memset(tp,0x8f,sizeof tp);
   if(v<=n){
    tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
    tp[1]=max(
     dp[u][0]+dp[v][0]+1
     dp[u][1]+max(dp[v][0],dp[v][1])
   }else{
    tp[0]=dp[u][0]+dp[v][0];
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
   dp[u][0]=tp[0],dp[u][1]=tp[1];
  }
 }else{
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
   dfs(v,u);
  min_dp[0][0]=0;
  min_dp[1][1]=1;
  min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
   memset(tmp,0x8f,sizeof tmp);
```

```
tmp[0][0]=max(
                                                                 ansd = d;
   min_dp[0][0]+max(dp[v][0],dp[v][1]),
                                                                 return:
    \min_{dp[0][1]+dp[v][0]}
                                                                int c = R[0];
                                                                for(int i = R[0]; i; i = R[i])
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
   tmp[1][0]=max(
                                                                 if(S[i] < S[c]) c = i;
   \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
                                                                remove(c);
                                                                for(int i = D[c]; i != c; i = D[i]) {
    min_dp[1][1]+dp[v][0]
                                                                 ans[d] = row[i];
                                                                 for(int j = R[i]; j != i; j = R[j])
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
                                                                  remove(col[j]);
   memcpy(min_dp,tmp,sizeof tmp);
                                                                 dance(d+1);
                                                                 for(int j = L[i]; j != i; j = L[j])
  dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
  dp[u][0]=min_dp[0][0];
                                                                  resume(col[j]);
                                                                resume(c);
int main(){
                                                               }
int m,a,b;
                                                              } sol;
scanf("%d%d",&n,&m);
for(int i=0;i<m;i++){
  scanf("%d%d",&a,&b);</pre>
                                                              8.8
                                                                    Tree Knapsack
                                                              int dp[N][K];PII obj[N];
  init_g[a].push_back(b);
                                                              vector<int> G[N];
 init_g[b].push_back(a);
                                                              void dfs(int u, int mx){
                                                               for(int s: G[u]) {
par[1]=-1;
                                                                if(mx < obj[s].first) continue;</pre>
tarjan(1);
                                                                for(int i=0;i<=mx-obj[s].FF;i++)</pre>
dfs(1,-1);
                                                                 dp[s][i] = dp[u][i];
printf("%d\n", max(dp[1][0], dp[1][1]));
                                                                dfs(s, mx - obj[s].first);
return 0;
                                                                for(int i=obj[s].FF;i<=mx;i++)</pre>
                                                                 dp[u][i] = max(dp[u][i],
                                                                  dp[s][i - obj[s].FF] + obj[s].SS);
8.7 DLX
struct DLX {
const static int maxn=210;
                                                              int main(){
const static int maxm=210;
                                                               int n, k; cin >> n >> k;
const static int maxnode=210*210;
                                                               for(int i=1;i<=n;i++){</pre>
int n, m, size, row[maxnode], col[maxnode];
                                                                int p; cin >> p;
int U[maxnode], D[maxnode], L[maxnode], R[maxnode];
                                                                G[p].push_back(i);
 int H[maxn], S[maxm], ansd, ans[maxn];
                                                                cin >> obj[i].FF >> obj[i].SS;
void init(int _n, int _m) {
 n = _n, m = _m;
                                                               dfs(0, k); int ans = 0;
 for(int i = 0; i <= m; ++i) {</pre>
                                                               for(int i=0;i<=k;i++) ans = max(ans, dp[0][i]);</pre>
  S[i] = 0;
                                                               cout << ans << '\n';
  U[i] = D[i] = i;
                                                               return 0:
  L[i] = i-1, R[i] = i+1;
 R[L[0] = size = m] = 0;
                                                              8.9 N Queens Problem
 for(int i = 1; i <= n; ++i) H[i] = -1;</pre>
                                                              vector< int > solve( int n ) {
                                                               // no solution when n=2, 3
void Link(int r, int c) {
                                                               vector< int > ret:
 ++S[col[++size] = c];
                                                               if ( n % 6 == 2 ) {
for ( int i = 2 ; i <= n ; i += 2 )
 row[size] = r; D[size] = D[c];
U[D[c]] = size; U[size] = c; D[c] = size;
                                                                 ret.push_back( i );
  if(H[r] < 0) H[r] = L[size] = R[size] = size;</pre>
                                                                ret.push_back( 3 ); ret.push_back( 1 );
 else {
                                                                for ( int i = 7 ; i <= n ; i += 2 )
  R[size] = R[H[r]];
                                                                 ret.push_back( i );
   L[R[H[r]]] = size;
                                                                ret.push_back( 5 );
  L[size] = H[r];
R[H[r]] = size;
                                                               } else if ( n % 6 == 3 ) {
for ( int i = 4 ; i <= n ; i += 2 )</pre>
                                                                 ret.push_back( i )
                                                                ret.push_back( 2 );
void remove(int c) {
                                                                for ( int i = 5 ; i <= n ; i += 2 )
 L[R[c]] = L[c]; R[L[c]] = R[c];
                                                                 ret.push_back( i );
  for(int i = D[c]; i != c; i = D[i])
                                                                ret.push_back( 1 ); ret.push_back( 3 );
   for(int j = R[i]; j != i; j = R[j]) {
                                                               } else {
   U[D[j]] = U[j];
                                                                for ( int i = 2 ; i <= n ; i += 2 )
   D[U[j]] = D[j];
                                                                 ret.push_back( i );
    --S[col[j]];
                                                                for ( int i = 1 ; i <= n ; i += 2 )
                                                                 ret.push_back( i );
void resume(int c)
 L[R[c]] = c; R[L[c]] = c;
                                                               return ret;
  for(int i = U[c]; i != c; i = U[i])
   for(int j = L[i]; j != i; j = L[j]) {
                                                              8.10 Aliens Optimization
   U[D[j]] = j;
   D[U[j]] = j;
                                                              long long Alien()
    ++S[col[j]];
                                                               long long c = kInf;
                                                               for (int d = 60; d >= 0; --d) {
   }
                                                                // cost can be negative, depending on the problem.
                                                                if (c - (1LL \ll d) < 0) continue;
void dance(int d) {
  if(d>=ansd) return;
                                                                long long ck = c - (1LL \ll d)
                                                                pair<long long, int> r = check(ck);
 if(R[0] == 0) {
```

```
if (r.second == k) return r.first - ck * k;
  if (r.second < k) c = ck;</pre>
pair<long long, int> r = check(c);
return r.first - c * k;
8.11 Hilbert Curve
long long hilbert(int n, int x, int y) {
long long res = 0;
for (int s = n / 2; s; s >>= 1) {
  int rx = (x & s) > 0, ry = (y & s) > 0;
  res += s * 111 * s * ((3 * rx) ^ ry);
 if (ry == 0) {
  if (rx == 1) x = s - 1 - x, y = s - 1 - y;
   swap(x, y);
  }
return res;
8.12 Binary Search On Fraction
struct Q {
11 p, q;
Q go(Q b, 11 d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 \le p,q \le N
Q frac_bs(11 N) {
Q lo{0, 1}, hi{1, 0};
if (pred(lo)) return lo;
 assert(pred(hi));
bool dir = 1, L = 1, H = 1;
for (; L || H; dir = !dir) {
  11 len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
   if (Q mid = hi.go(lo, len + step);
     mid.p > N || mid.q > N || dir ^ pred(mid))
   t++;
else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
return dir ? hi : lo;
```