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1
 Basic
```

1.1 vimrc

```
se is nu bs=2 ru mouse=a encoding=utf-8 ls=2
se cin cino+=j1 et sw=4 sts=4 tgc sc hls
syn on
colorscheme desert
filetype indent on
inoremap \{ < CR > \{ < CR > \} < ESC > 0 \}
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -std=c++17 -
    DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
    Wconversion -fsanitize=address,undefined -g && echo
     success<CR>
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 &&
     echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
```

1.2 Debug Macro

```
#ifdef KISEKI
#define safe cerr<<__PRETTY_FUNCTION__\
<<" line "<<__LINE__<<" safe\n"</pre>
#define debug(a...) qwerty(#a, a)
#define orange(a...) dvorak(#a, a)
using std::cerr;
template <typename ...T>
void qwerty(const char *s, T ...a) {
  cerr << "\e[1;32m(" << s << ") = (";
  int cnt = sizeof...(T);</pre>
   (..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
}
template <typename Iter>
void dvorak(const char *s, Iter L, Iter R) {
  cerr << "\e[1;32m[ " << s << " ] = [ ";</pre>
  for (int f = 0; L != R; ++L)
    cerr << (f++? ", ": "") << *L;
    cerr << "]\e[0m\n";
}
#else
#define safe ((void)0)
#define debug(...) ((void)0)
#define orange(...) ((void)0)
#endif
```

1.3 Increase Stack

```
const int size = 256 << 20;</pre>
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));
```

1.4 Pragma Optimization

```
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
1.5 IO Optimization
static inline int gc() {
  constexpr int B = 1 << 20;
  static char buf[B], *p, *q;
  if(p == a \&\&
    (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
   return EOF;
  return *p++;
template < typename T >
static inline bool gn( T &x ) {
 int c = gc(); T sgn = 1; x = 0;
while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
if(c == '-') sgn = -1, c = gc();
 if(c == EOF) return false;
 while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
 return x *= sgn, true;
```

#pragma GCC optimize("Ofast,no-stack-protector")

#pragma GCC optimize("no-math-errno,unroll-loops")

2 Data Structure

2.1 Dark Magic

2.2 Link-Cut Tree

p->ch[dir]=c;

```
struct Node{
Node *par, *ch[2];
int xor_sum, v;
bool is_rev;
Node(int _v){
 v=xor_sum=_v;is_rev=false;
 par=ch[0]=ch[1]=nullptr;
inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
inline void down(){
 if(is_rev){
  if(ch[0]!=nullptr) ch[0]->set_rev();
   if(ch[1]!=nullptr) ch[1]->set_rev();
   is_rev=false;
 }
inline void up(){
 xor_sum=v;
  if(ch[0]!=nullptr){
  xor_sum^=ch[0]->xor_sum;
  ch[0]->par=this;
 if(ch[1]!=nullptr){
  xor_sum^=ch[1]->xor_sum;
  ch[1]->par=this;
inline bool is_root(){
 return par==nullptr ||\
   (par->ch[0]!=this && par->ch[1]!=this);
bool is_rch(){return !is_root() && par->ch[1]==this;}
} *node[maxn], *stk[maxn];
int top;
void to_child(Node* p,Node* c,bool dir){
```

```
p->up();
inline void rotate(Node* node){
 Node* par=node->par;
 Node* par_par=par->par;
 bool dir=node->is_rch()
 bool par_dir=par->is_rch()
 to_child(par, node->ch[!dir], dir);
 to_child(node,par,!dir);
 if(par_par!=nullptr && par_par->ch[par_dir]==par)
  to_child(par_par,node,par_dir);
 else node->par=par_par;
inline void splay(Node* node){
 Node* tmp=node;
 stk[top++]=node;
 while(!tmp->is_root()){
  tmp=tmp->par;
  stk[top++]=tmp;
 while(top) stk[--top]->down();
 for(Node *fa=node->par;
  !node->is_root();
  rotate(node), fa=node->par)
  if(!fa->is_root())
   rotate(fa->is_rch()==node->is_rch()?fa:node);
inline void access(Node* node){
 Node* last=nullptr;
 while(node!=nullptr){
  splay(node);
  to_child(node, last, true);
  last=node;
  node=node->par;
inline void change_root(Node* node){
 access(node);splay(node);node->set_rev();
inline void link(Node* x, Node* y){
 change_root(x);splay(x);x->par=y;
inline void split(Node* x,Node* y){
 change_root(x);access(y);splay(x);
 to_child(x,nullptr,true);y->par=nullptr;
inline void change_val(Node* node,int v){
access(node);splay(node);node->v=v;node->up();
inline int query(Node* x,Node* y){
 change_root(x);access(y);splay(y);
 return y->xor_sum;
inline Node* find_root(Node* node){
 access(node);splay(node);
 Node* last=nullptr:
 while(node!=nullptr){
  node->down();last=node;node=node->ch[0];
 return last;
set<pii> dic;
inline void add_edge(int u,int v){
 if(u>v) swap(u,v)
 if(find_root(node[u])==find_root(node[v])) return;
 dic.insert(pii(u,v))
link(node[u],node[v]);
inline void del_edge(int u,int v){
 if(u>v) swap(u,v);
 if(dic.find(pii(u,v))==dic.end()) return;
 dic.erase(pii(u,v))
 split(node[u],node[v]);
2.3 LiChao Segment Tree
struct Line{
 int m, k, id;
 Line() : id( -1 ) {}
Line('int a, int'b,'int c')
: m(a), k(b), id(c) {}
```

int at(int x) { return m * x + k; }

```
#undef sz
class LiChao {
 private:
                                                             2.5 Linear Basis
  int n; vector< Line > nodes;
  inline int lc( int x ) { return 2 * x + 1; }
                                                             template <int BITS>
  inline int rc( int x ) { return 2 * x + 2; }
                                                             struct LinearBasis {
  void insert( int 1, int r, int id, Line ln ) {
                                                              array<uint64_t, BITS> basis;
   int m = (1 + r) >> 1;
                                                              Basis() { basis.fill(0); }
                                                              void add(uint64_t x)
   if ( nodes[ id ].id == -1 ) {
    nodes[ id ] = ln;
                                                               for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
                                                                if (basis[i] == 0) {
    return:
                                                                 basis[i] = x;
   bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
                                                                 return;
   if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
    atLeft ^= 1; swap( nodes[ id ], ln );
                                                                x ^= basis[i];
                                                               }
   if ( r - 1 == 1 ) return;
   if ( atLeft ) insert( l, m, lc( id ), ln );
                                                              bool ok(uint64_t x) {
   else insert( m, r, rc( id ), ln );
                                                               for (int i = 0; i < BITS; ++i)</pre>
                                                                if ((x >> i) & 1) x ^= basis[i];
  int query( int 1, int r, int id, int x ) {
                                                               return x == 0;
   int ret = 0;
   if ( nodes[ id ].id != -1 )
                                                             };
    ret = nodes[ id ].at( x );
                                                             2.6
                                                                   Binary Search On Segment Tree
   int m = (1 + r) >> 1;
   if ( r - l == 1 ) return ret;
                                                             // find_first = x -> minimal x s.t. check( [a, x) )
   else if (x < m )
                                                             // find_last = x \rightarrow maximal x s.t. check([x, b))
    return max( ret, query( 1, m, lc( id ), x ) );
                                                             template <typename C>
   else
                                                             int find_first(int 1, const C &check) {
    return max( ret, query( m, r, rc( id ), x ) );
                                                              if (1 >= n) return n;
                                                              1 += sz;
 public:
                                                              for (int i = height; i > 0; i--)
  void build( int n_ ) {
                                                               propagate(l >> i);
  n = n_; nodes.clear();
                                                              Monoid sum = identity;
   nodes.resize( n << 2, Line() );</pre>
                                                               while ((1 & 1) == 0) 1 >>= 1;
  void insert( Line ln ) { insert( 0, n, 0, ln ); }
                                                               if (check(f(sum, data[1]))) {
  int query( int x ) { return query( 0, n, 0, x ); }
                                                                while (1 < sz) {</pre>
                                                                 propagate(1);
                                                                 1 <<= 1;
2.4 Treap
                                                                 auto nxt = f(sum, data[1]);
namespace Treap{
                                                                 if (not check(nxt)) {
 #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                                  sum = nxt;
 struct node{
                                                                  1++;
  int size;
                                                                 }
  uint32_t pri;
                                                                }
  node *lc, *rc, *pa;
                                                                return 1 + 1 - sz;
  node():size(0),pri(rand()),lc(0),rc(0),pa(0){}
  void pull() {
                                                               sum = f(sum, data[1++]);
  size = 1; pa = nullptr;
                                                              } while ((1 & -1) != 1);
   if ( lc ) { size += lc->size; lc->pa = this; }
if ( rc ) { size += rc->size; rc->pa = this; }
                                                              return n;
  }
                                                             template <typename C>
                                                             int find_last(int r, const C &check) {
node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
                                                              if (r <= 0) return -1;
                                                              r += sz;
  if ( L->pri > R->pri ) {
                                                              for (int i = height; i > 0; i--)
  L->rc = merge( L->rc, R ); L->pull();
                                                               propagate((r - 1) >> i);
   return L;
                                                              Monoid sum = identity;
  } else {
                                                              do {
   R->lc = merge( L, R->lc ); R->pull();
   return R;
                                                               while (r > 1 \text{ and } (r \& 1)) r >>= 1;
  }
                                                               if (check(f(data[r], sum))) {
                                                                while (r < sz) {</pre>
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                                 propagate(r);
  if ( not rt ) L = R = nullptr;
                                                                 r = (r << 1) + 1;
  else if( sz( rt->lc ) + 1 <= k ) {
                                                                 auto nxt = f(data[r], sum);
                                                                 if (not check(nxt)) {
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                  sum = nxt;
   L->pull();
                                                                  r--:
  } else {
                                                                 }
   R = rt;
   split_by_size( rt->lc, k, L, R->lc );
                                                                return r - sz;
   R->pull();
                                                               sum = f(data[r], sum);
  }
                                                              } while ((r & -r) != r);
 int getRank(node *o) {
                                                              return -1;
  int r = sz(o->lc);
  for (;o->pa != nullptr; o = o->pa)
   if (o->pa->rc != o) r += sz(o->pa->lc);
                                                                  Graph
  return r;
```

3.1 BCC Edge

```
class BCC_Bridge {
                                                                  for (int i = 0; i < n; ++i)
                                                                   if (not dfn[i]) dfs(i, i);
private:
 int n, ecnt;
  vector<vector<pair<int,int>>> G;
                                                                 int get_id(int x) { return bcc[x]; }
 vector<int> dfn, low;
                                                                 int count() { return ecnt; }
  vector<bool> bridge;
                                                                 bool is_ap(int x) { return ap[x]; }
 void dfs(int u, int f) {
   dfn[u] = low[u] = dfn[f] + 1;
                                                              } bcc_ap;
                                                               3.3 2-SAT (SCC)
   for (auto [v, t]: G[u]) {
  if (v == f) continue;
                                                              class TwoSat{
    if (dfn[v]) {
                                                                private:
     low[u] = min(low[u], dfn[v]);
                                                                 int n:
     continue;
                                                                 vector<vector<int>> rG,G,sccs;
                                                                 vector<int> ord,idx;
    dfs(v, u);
                                                                 vector<bool> vis,result;
   low[u] = min(low[u], low[v]);
if (low[v] > dfn[u]) bridge[t] = true;
                                                                 void dfs(int u){
                                                                  vis[u]=true
                                                                  for(int v:G[u])
  }
                                                                   if(!vis[v]) dfs(v);
public:
                                                                  ord.push_back(u);
  void init(int n_) {
   G.clear(); G.resize(n = n_);
                                                                 void rdfs(int u){
   low.assign(n, ecnt = 0);
                                                                  vis[u]=false;idx[u]=sccs.size()-1;
                                                                  sccs.back().push_back(u);
   dfn.assign(n, 0);
                                                                  for(int v:rG[u])
  void add_edge(int u, int v) {
                                                                   if(vis[v])rdfs(v);
  G[u].emplace_back(v, ecnt);
   G[v].emplace_back(u, ecnt++);
                                                                public:
                                                                 void init(int n_){
  void solve() {
                                                                  n=n_;G.clear();G.resize(n);
  bridge.assign(ecnt, false);
                                                                  rG.clear();rG.resize(n)
   for (int i = 0; i < n; ++i)
                                                                  sccs.clear();ord.clear();
    if (not dfn[i]) dfs(i, i);
                                                                  idx.resize(n);result.resize(n);
  bool is_bridge(int x) { return bridge[x]; }
                                                                 void add_edge(int u,int v){
} bcc_bridge;
                                                                  G[u].push_back(v);rG[v].push_back(u);
3.2 BCC Vertex
                                                                 void orr(int x,int y){
class BCC_AP {
                                                                  if ((x^y)==1)return
                                                                  add_edge(x^1,y); add_edge(y^1,x);
private:
 int n, ecnt;
 vector<vector<pair<int,int>>> G;
                                                                 bool solve(){
  vector<int> bcc, dfn, low, st;
                                                                  vis.clear();vis.resize(n);
  vector<bool> ap, ins;
                                                                  for(int i=0;i<n;++i)</pre>
 void dfs(int u, int f) {
  dfn[u] = low[u] = dfn[f] + 1;
                                                                   if(not vis[i])dfs(i);
                                                                  reverse(ord.begin(),ord.end());
   int ch = 0;
                                                                  for (int u:ord){
   for (auto [v, t]: G[u]) if (v != f) {
                                                                   if(!vis[u])continue:
    if (not ins[t]) {
                                                                   sccs.push_back(vector<int>());
     st.push_back(t);
                                                                   rdfs(u);
     ins[t] = true;
                                                                  for(int i=0;i<n;i+=2)</pre>
    if (dfn[v]) {
                                                                   if(idx[i]==idx[i+1])
     low[u] = min(low[u], dfn[v]);
                                                                    return false;
                                                                  vector<bool> c(sccs.size());
    } ++ch; dfs(v, u);
                                                                  for(size_t i=0;i<sccs.size();++i){</pre>
    low[u] = min(low[u], low[v]);
                                                                   for(size_t j=0;j<sccs[i].size();++j){</pre>
                                                                    result[sccs[i][j]]=c[i]
    if (low[v] >= dfn[u]) {
     ap[u] = true;
                                                                    c[idx[sccs[i][j]^1]]=!c[i];
     while (true) {
      int eid = st.back(); st.pop_back();
      bcc[eid] = ecnt;
                                                                  return true;
      if (eid == t) break;
     }
                                                                 bool get(int x){return result[x];}
                                                                 inline int get_id(int x){return idx[x];}
     ecnt++;
                                                                 inline int count(){return sccs.size();}
    }
                                                              } sat2;
   if (ch == 1 and u == f) ap[u] = false;
                                                               3.4 Lowbit Decomposition
public:
                                                              class LowbitDecomp{
  void init(int n_) {
                                                               private:
  G.clear(); G.resize(n = n_);
                                                                int time_, chain_, LOG_N;
   ecnt = 0; ap.assign(n, false);
                                                                vector< vector< int > > G, fa;
                                                                vector< int > tl, tr, chain, chain_st;
// chain_ : number of chain
   low.assign(n, 0); dfn.assign(n, 0);
                                                                // tl, tr[ u ] : subtree interval in the seq. of u
  void add_edge(int u, int v) {
                                                                // chain_st[ u ] : head of the chain contains u // chian[ u ] : chain id of the chain u is on
   G[u].emplace_back(v, ecnt);
   G[v].emplace_back(u, ecnt++);
                                                                void predfs( int u, int f ) {
  void solve() {
                                                                 chain[ u ] = 0;
   ins.assign(ecnt, false);
                                                                 for ( int v : G[ u ] ) {
                                                                  if ( v == f ) continue;
   bcc.resize(ecnt); ecnt = 0;
```

```
predfs( v, u );
                                                              class MaxClique{
   if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
                                                              private:
    chain[ u ] = chain[ v ];
                                                               using bits = bitset< MAXN >;
                                                               bits popped, G[ MAXN ], ans;
size_t deg[ MAXN ], deo[ MAXN ], n;
  if ( not chain[ u ] )
   chain[ u ] = chain_ ++;
                                                               void sort_by_degree() {
                                                                popped.reset();
                                                                for ( size_t i = 0 ; i < n ; ++ i )</pre>
 void dfschain( int u, int f ) {
  fa[ u ][ 0 ] = f;
for ( int i = 1 ; i < LOG_N ; ++ i )
                                                                  deg[ i ] = G[ i ].count();
                                                                for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;</pre>
   fa[u][i] = fa[fa[u][i-1]][i-1];
                                                                  for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )
    mi = deg[ id = j ];</pre>
  tl[ u ] = time_++;
  if ( not chain_st[ chain[ u ] ] )
   chain_st[ chain[ u ] ] = u;
                                                                  popped[ deo[ i ] = id ] = 1;
  for ( int v : G[ u ] )
   if ( v != f and chain[ v ] == chain[ u ] )
                                                                  for( size_t u = G[ i ]._Find_first() ;
  dfschain( v, u );
for ( int v : G[ u ] )
                                                                   u < n ; u = G[ i ]._Find_next( u ) )</pre>
                                                                     -- deg[ u ];
   if ( v != f and chain[ v ] != chain[ u ] )
                                                                }
    dfschain( v, u );
                                                               void BK( bits R, bits P, bits X ) {
  tr[ u ] = time_;
                                                                if (R.count()+P.count() <= ans.count()) return;</pre>
                                                                if ( not P.count() and not X.count() ) {
 bool anc( int u, int v )
 return tl[ u ] <= tl[ v ] and tr[ v ] <= tr[ u ];</pre>
                                                                 if ( R.count() > ans.count() ) ans = R;
                                                                 return:
public:
                                                                }
                                                                /* greedily chosse max degree as pivot
 int lca( int u, int v ) {
  if ( anc( u, v ) ) return u;
                                                                bits cur = P | X; size_t pivot = 0, sz = 0;
  for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
                                                                for ( size_t u = cur._Find_first() ;
   if ( not anc( fa[ u ][ i ], v ) )
                                                                 u < n ; u = cur._Find_next( u )
    u = fa[ u ][ i ];
                                                                  if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
                                                                cur = P & ( ~G[ pivot ] );
  return fa[ u ][ 0 ];
                                                                 */ // or simply choose first
                                                                bits cur = P & (~G[ ( P | X )._Find_first() ]);
 void init( int n ) {
  fa.assign( ++n, vector< int > ( LOG_N ) );
                                                                for ( size_t u = cur._Find_first()
  for (LOG_N = 0 ; (1 << LOG_N ) < n ; ++ LOG_N );
                                                                 u < n ; u = cur._Find_next( u ) ) {
                                                                 if ( R[ u ] ) continue;
  G.clear(); G.resize( n );
  tl.assign( n, 0 ); tr.assign( n, 0 )
                                                                 R[u] = 1;
  chain.assig( n, 0 ); chain_st.assign( n, 0 );
                                                                 BK( R, P & G[ u ], X & G[ u ] );
                                                                 R[u] = P[u] = 0, X[u] = 1;
 void add_edge( int u , int v ) {
  // 1-base
  G[ u ].push_back( v );
                                                              public:
  G[ v ].push_back( u );
                                                               void init( size_t n_ ) {
                                                                n = n_{-};
 }
 void decompose(){
                                                                for ( size_t i = 0 ; i < n ; ++ i )
                                                                 G[ i ].reset();
 chain_ = 1;
 predfs( 1, 1 );
                                                                ans.reset();
  time_{-} = 0;
 dfschain( 1, 1 );
                                                               void add_edges( int u, bits S ) { G[ u ] = S; }
                                                               void add_edge( int u, int v ) {
 PII get_subtree(int u) { return {tl[ u ],tr[ u ] }; }
                                                                G[u][v] = G[v][u] = 1;
 vector< PII > get_path( int u , int v ){
  vector< PII > res;
                                                               int solve() {
  int g = lca( u, v );
                                                                sort_by_degree(); // or simply iota( deo... )
  while ( chain[ u ] != chain[ g ] ) {
                                                                for ( size_t i = 0 ; i < n ; ++ i )</pre>
   int s = chain_st[ chain[ u ] ];
                                                                 deg[ i ] = G[ i ].count();
   res.emplace_back( tl[ s ], tl[ u ] + 1 );
                                                                bits pob, nob = 0; pob.set();
   u = fa[ s ][ 0 ];
                                                                for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
                                                                for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
  res.emplace_back( tl[ g ], tl[ u ] + 1 );
while ( chain[ v ] != chain[ g ] ) {
                                                                 size_t v = deo[ i ];
                                                                 bits tmp; tmp[ v ] = 1;
  int s = chain_st[ chain[ v ] ];
                                                                 BK( tmp, pob & G[ v ], nob & G[ v ] );
   res.emplace_back( tl[ s ], tl[ v ] + 1 );
                                                                 pob[v] = 0, nob[v] = 1;
   v = fa[ s ][ 0 ];
                                                                return static_cast< int >( ans.count() );
  res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
                                                               }
  return res;
                                                              };
  /* res : list of intervals from u to v
                                                                   MaxCliqueDyn
   \star ( note only nodes work, not edge )
                                                              constexpr int kN = 150;
   * vector< PII >& path = tree.get_path( u , v )
                                                              struct MaxClique { // Maximum Clique
                                                               bitset<kN> a[kN], cs[kN];
   * for( auto [ 1, r ] : path ) {
   * 0-base [ 1, r )
                                                               int ans, sol[kN], q, cur[kN], d[kN], n;
   * }
                                                               void init(int _n) {
   */
                                                               n = n, ans q = 0;
                                                                for (int i = 0; i < n; i++) a[i].reset();</pre>
} tree;
                                                               void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
3.5 MaxClique
                                                               void csort(vector<int> &r, vector<int> &c) {
// contain a self loop u to u, than u won't in clique
                                                                int mx = 1, km = max(ans - q + 1, 1), t = 0,
                                                                  m = int(r.size());
template < size_t MAXN >
```

```
cs[1].reset(); cs[2].reset();
                                                                 addEdge(stk[i], stk[i + 1]);
  for (int i = 0; i < m; i++) {
  int p = r[i], k = 1;</pre>
                                                               3.8 Centroid Decomposition
   while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
                                                               struct Centroid {
   cs[k][p] = 1;
                                                                vector<vector<int64_t>> Dist;
   if (k < km) r[t++] = p;
                                                                vector<int> Parent, Depth;
                                                                vector<int64_t> Sub, Sub2;
  c.resize(m);
                                                                vector<int> Sz, Sz2;
  if (t) c[t - 1] = 0;
                                                                Centroid(vector<vector<pair<int, int>>> g) {
  for (int k = km; k <= mx; k++) {</pre>
                                                                 int N = g.size()
  for (int p = int(cs[k]._Find_first());
                                                                 vector<bool> Vis(N);
      p < kN; p = int(cs[k]._Find_next(p))) {</pre>
                                                                 vector<int> sz(N), mx(N);
    r[t] = p; c[t++] = k;
                                                                 vector<int> Path;
                                                                 Dist.resize(N)
                                                                 Parent.resize(N);
                                                                 Depth.resize(N)
 void dfs(vector<int> &r, vector<int> &c, int 1,
                                                                 auto DfsSz = [\&](auto dfs, int x) -> void {
  bitset<kN> mask) {
                                                                  Vis[x] = true; sz[x] = 1; mx[x] = 0;
  while (!r.empty()) {
                                                                  for (auto [u, w] : g[x]) {
                                                                   if (Vis[u]) continue;
   int p = r.back(); r.pop_back();
   mask[p] = 0;
                                                                   dfs(dfs, u)
   if (q + c.back() <= ans) return;</pre>
                                                                   sz[x] += sz[u];
   cur[q++] = p;
                                                                   mx[x] = max(mx[x], sz[u]);
   vector<int> nr, nc;
   bitset<kN> nmask = mask & a[p];
                                                                  Path.push_back(x);
   for (int i : r)
                                                                 };
    if (a[p][i]) nr.push_back(i);
                                                                 auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
   if (!nr.empty()) {
                                                                  -> void {
                                                                  Dist[x].push_back(D);Vis[x] = true;
    if (1 < 4) {
     for (int i : nr)
                                                                  for (auto [u, w] : g[x]) {
                                                                   if (Vis[u]) continue;
      d[i] = int((a[i] & nmask).count());
     sort(nr.begin(), nr.end(),
                                                                   dfs(dfs, u, D + w);
      [&](int x, int y)
       return d[x] > d[y];
      });
                                                                 auto Dfs = [&]
                                                                  (auto dfs, int x, int D = 0, int p = -1)->void {
   csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
                                                                  Path.clear(); DfsSz(DfsSz, x);
                                                                  int M = Path.size();
    ans = q; copy(cur, cur + q, sol);
                                                                  int C = -1:
                                                                  for (int u : Path) {
   c.pop_back(); q--;
                                                                   if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
  }
                                                                   Vis[u] = false;
 int solve(bitset<kN> mask) { // vertex mask
                                                                  DfsDist(DfsDist, C);
                                                                  for (int u : Path) Vis[u] = false;
  vector<int> r, c;
  for (int i = 0; i < n; i++)
                                                                  Parent[C] = p; Vis[C] = true;
  if (mask[i]) r.push_back(i);
for (int i = 0; i < n; i++)</pre>
                                                                  Depth[C] = D;
                                                                  for (auto [u, w] : g[C]) {
  d[i] = int((a[i] & mask).count());
                                                                   if (Vis[u]) continue;
  sort(r.begin(), r.end(),
                                                                   dfs(dfs, u, D + 1, C);
   [&](int i, int j) { return d[i] > d[j]; });
  csort(r, c);
  dfs(r, c, 1, mask);
                                                                 Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
  return ans; // sol[0 ~ ans-1]
                                                                 Sz.resize(N); Sz2.resize(N);
} graph;
                                                                void Mark(int v) {
                                                                 int x = v, z = -1
3.7 Virtural Tree
                                                                 for (int i = Depth[v]; i >= 0; --i) {
                                                                  Sub[x] += Dist[v][i]; Sz[x]++;
inline bool cmp(const int &i, const int &j) {
                                                                  if (z != -1) {
return dfn[i] < dfn[j];</pre>
                                                                   Sub2[z] += Dist[v][i];
void build(int vectrices[], int k) {
                                                                   Sz2[z]++;
 static int stk[MAX_N];
 sort(vectrices, vectrices + k, cmp);
                                                                  z = x; x = Parent[x];
 stk[sz++] = 0;
                                                                 }
 for (int i = 0; i < k; ++i) {
  int u = vectrices[i], lca = LCA(u, stk[sz - 1]);
  if (lca == stk[sz - 1]) stk[sz++] = u;</pre>
                                                                int64_t Query(int v) {
                                                                 int64_t res = 0;
                                                                 int x = v, z = -1
                                                                 for (int i = Depth[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
   while (sz \ge 2 \&\& dep[stk[sz - 2]] \ge dep[lca]) {
    addEdge(stk[sz - 2], stk[sz - 1]);
                                                                  if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
                                                                  z = x; x = Parent[x];
   if (stk[sz - 1] != lca) {
   addEdge(lca, stk[--sz]);
                                                                 return res;
    stk[sz++] = lca, vectrices[cnt++] = lca;
                                                               };
   stk[sz++] = u;
                                                               3.9 Tree Hashing
  }
                                                              |uint64_t hsah(int u, int f) {
 for (int i = 0; i < sz - 1; ++i)
                                                              uint64_t r = 127;
```

```
for (int v : G[ u ]) if (v != f) {
  uint64_t hh = hsah(v, u);
                                                                        stk[ stk_ ++ ] = u;
  r=(r+(hh*hh)%1010101333)%1011820613;
                                                                       bool inPath[ N ];
                                                                       void Diff( int u ) {
return r:
                                                                        if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
}
                                                                        else { /*add this edge*/ }
3.10 Minimum Mean Cycle
/* minimum mean cycle O(VE) */
                                                                       void traverse( int& origin_u, int u ) {
                                                                        for ( int g = lca( origin_u, u )
struct MMC{
                                                                         origin_u != g ; origin_u = parent_of[ origin_u ] )
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
                                                                          Diff( origin_u );
#define V 1021
                                                                        for (int v = u; v != origin_u; v = parent_of[v])
#define inf 1e9
                                                                         Diff( v );
 struct Edge { int v,u; double c; };
                                                                        origin_u = u;
 int n, m, prv[V][V], prve[V][V], vst[V];
                                                                       void solve() {
 Edge e[E];
                                                                        dfs( 1, 1 );
 vector<int> edgeID, cycle, rho;
 double d[V][V];
                                                                        while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
                                                                        sort( que, que + q, [](const Que& x, const Que& y) {
  return tie( block_id[ x.u ], dfn[ x.v ] )
 void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
 void add_edge( int vi , int ui , double ci )
                                                                              < tie( block_id[ y.u ], dfn[ y.v ] );
 { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
                                                                        } );
                                                                        int U = 1, V = 1;
  for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);</pre>
                                                                        for ( int i = 0 ; i < q ; ++ i ) {
  pass( U, que[ i ].u );
  pass( V, que[ i ].v );</pre>
   for(int j=0; j<m; j++) {</pre>
                                                                         // we could get our answer of que[ i ].id
    int v = e[j].v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
                                                                      }
                                                                       /*
     d[i+1][u] = d[i][v]+e[j].c;
      prv[i+1][u] = v;
                                                                      Method 2:
      prve[i+1][u] = j;
                                                                      dfs u:
                                                                       push u
                                                                        iterate subtree
                                                                       Let P = LCA(u, v), and St(u) \le St(v)
                                                                      if (P == u) query[St(u), St(v)]
 double solve(){
  // returns inf if no cycle, mmc otherwise
                                                                       else query[Ed(u), St(v)], query[St(P), St(P)]
  double mmc=inf;
  int st = -1;
                                                                       3.12 Minimum Steiner Tree
  bellman_ford();
  for(int i=0; i<n; i++) {</pre>
                                                                      // Minimum Steiner Tree
   double avg=-inf;
                                                                       // 0(V 3^T + V^2 2^T)
                                                                      struct SteinerTree{
   for(int k=0; k<n; k++) {</pre>
    if(d[n][i]<inf-eps)</pre>
                                                                       #define V 33
      avg=max(avg,(d[n][i]-d[k][i])/(n-k));
                                                                       #define T 8
                                                                       #define INF 1023456789
     else avg=max(avg,inf);
                                                                        int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
                                                                        void init( int _n ){
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
                                                                         n = _n;
                                                                         for( int i = 0 ; i < n ; i ++ ){
  for( int j = 0 ; j < n ; j ++ )
  dst[ i ][ j ] = INF;</pre>
  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
  for (int i=n; !vst[st]; st=prv[i--][st]) {
   vst[st]++
   edgeID.PB(prve[i][st]);
                                                                          dst[ i ][ i ] = 0;
   rho.PB(st);
  while (vst[st] != 2) {
  int v = rho.back(); rho.pop_back();
                                                                        void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
   cycle.PB(v);
                                                                         dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
   vst[v]++;
                                                                        void shortest_path(){
  reverse(ALL(edgeID));
                                                                         for( int k = 0 ; k < n ; k ++ )</pre>
                                                                          for( int i = 0 ; i < n ; i ++ )
for( int j = 0 ; j < n ; j ++ )
dst[ i ][ j ] = min( dst[ i ][ j ],
  edgeID.resize(SZ(cycle));
  return mmc;
} mmc;
                                                                                 dst[ i ][ k ] + dst[ k ][ j ] );
3.11 Mo's Algorithm on Tree
                                                                        int solve( const vector<int>& ter ){
                                                                         int t = (int)ter.size();
int q; vector< int > G[N];
                                                                         for( int i = 0 ; i < (1 << t) ; i ++ )
for( int j = 0 ; j < n ; j ++ )
struct Que{
 int u, v, id;
                                                                            dp[ i ][ j ] = INF;
} que[ N ];
int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
void dfs( int u, int f ) {
                                                                         for( int i = 0 ; i < n ; i ++ )</pre>
                                                                          dp[0][i] = 0;
 dfn[ u ] = dfn_++; int saved_rbp = stk_;
                                                                         for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
 for ( int v : G[ u ] ) {
  if ( v == f ) continue;
                                                                          if( msk == ( msk & (-msk) ) ){
                                                                           int who = __lg( msk );
  dfs( v, u );
                                                                            for( int i = 0 ; i < n ; i ++ )</pre>
  if ( stk_ - saved_rbp < SQRT_N ) continue;
for ( ++ block_ ; stk_ != saved_rbp ; )
  block_id[ stk[ -- stk_ ] ] = block_;</pre>
                                                                             dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
                                                                            continue;
```

```
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   for( int i = 0 ; i < n ; i ++ )</pre>
    for( int submsk = ( msk - 1 ) & msk ; submsk ;
          submsk = ( submsk - 1 ) & msk )
       dp[ msk ^ submsk ][ i ] );
   for( int i = 0 ; i < n ; i ++ ){</pre>
    tdst[ i ] = INF;
    for( int j = 0 ; j < n ; j ++ )
tdst[ i ] = min( tdst[ i ],
    dp[ msk ][ j ] + dst[ j ][ i ] );</pre>
   for( int i = 0 ; i < n ; i ++ )
dp[ msk ][ i ] = tdst[ i ];</pre>
  int ans = INF;
  for( int i = 0 ; i < n ; i ++ )
ans = min( ans , dp[ ( 1 << t ) - 1 ][ i ] );
  return ans;
} solver;
      Directed Minimum Spanning Tree
struct DirectedMST { // find maximum
 struct Edge {
  int u, v;
  int w;
  Edge(int u, int v, int w) : u(u), v(v), w(w) {}
 vector<Edge> Edges;
 void clear() { Edges.clear(); }
```

```
void addEdge(int a, int b, int w) { Edges.emplace_back
 (a, b, w); }
int solve(int root, int n) {
  vector<Edge> E = Edges;
  int ans = 0;
  while (true) {
   // find best in edge
   vector<int> in(n, -inf), prv(n, -1);
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
     in[e.v] = e.w;
     prv[e.v] = e.u;
   in[root] = 0;
   prv[root] = -1;
   for (int i = 0; i < n; i++)
    if (in[i] == -inf)
     return -inf;
   // find cycle
   int tot = 0;
   vector<int> id(n, -1), vis(n, -1);
for (int i = 0; i < n; i++) {</pre>
    ans += in[i];
    for (int x = i; x != -1 && id[x] == -1; x = prv[x])
     if (vis[x] == i) {
      for (int y = prv[x]; y != x; y = prv[y])
       id[y] = tot;
      id[x] = tot++;
      break;
     vis[x] = i;
    }
   if (!tot)
    return ans;
   for (int i = 0; i < n; i++)</pre>
    if (id[i] == -1)
     id[i] = tot++;
   // shrink
   for (auto &e : E) {
    if (id[e.u] != id[e.v])
     e.w -= in[e.v];
    e.u = id[e.u], e.v = id[e.v];
   n = tot:
   root = id[root];
  assert(false);
} DMST;
```

3.14 Manhattan Minimum Spanning Tree

```
typedef Point<int> P;
vector<array<int, 3>> manhattanMST(vector<P> ps) {
 vi id(sz(ps));
 iota(all(id), 0);
 vector<array<int, 3>> edges;
 rep(k, 0, 4) {
  sort(all(id), [&](int i, int j) {
   return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
  });
  map<int, int> sweep;
  for (int i : id) {
   for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++)) {
    int j = it->second;
    P d = ps[i] - ps[j];
    if (d.y > d.x) break;
    edges.push_back(\{d.y + d.x, i, j\});
   sweep[-ps[i].y] = i;
  for (P &p : ps)
   if (k \& 1) p.x = -p.x;
   else swap(p.x, p.y);
 return edges; // [{w, i, j}, ...]
}
```

```
Dominator Tree
3.15
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
 // vertices are numbered from 0 to n-1
 fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
 fill(fa, fa + n, -1); fill(val, val + n, -1);
 fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
 fill(dom, dom + n, -1); tk = 0;
 for (int i = 0; i < n; ++i) {
  g[i].clear(); r[i].clear(); rdom[i].clear();
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
 rev[dfn[x] = tk] = x;
 fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
 for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
  r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
 if (fa[x] == x) return c ? -1 : x;
 int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
 if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
 fa[x] = p;
 return c ? p : val[x];
vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
// p[i] = -2 if i is unreachable from s
 dfs(s);
 for (int i = tk - 1; i >= 0; --i) {
  for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
  if (i) rdom[sdom[i]].push_back(i);
  for (int &u : rdom[i]) {
   int p = find(u);
   if (sdom[p] == i) dom[u] = i;
   else dom[u] = p;
  if (i) merge(i, rp[i]);
 vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)</pre>
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
 for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
 return p;
}}
```

3.16 Edge Coloring

```
// \max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
for (int i = 0; i <= N; i++)
 for (int j = 0; j <= N; j++)
C[i][j] = G[i][j] = 0;</pre>
void solve(vector<pair<int, int>> &E, int N) {
int X[kN] = {}, a;
auto update = [&](int u)
 for (X[u] = 1; C[u][X[u]]; X[u]++);
auto color = [&](int u, int v, int c) {
 int p = G[u][v];
G[u][v] = G[v][u] = c;
 C[u][c] = v, C[v][c] = u;
 C[u][p] = C[v][p] = 0;
  if(p) X[u] = X[v] = p
 else update(u), update(v);
 return p;
 };
auto flip = [&](int u, int c1, int c2) {
 int p = C[u][c1];
 swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
 if (!C[u][c1]) X[u] = c1;
 if (!C[u][c2]) X[u] = c2;
  return p;
for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {
 auto [u, v] = E[t];
  int v0 = v, c = X[u], c0 = c, d;
  vector<pair<int,
                    int>> L; int vst[kN] = {};
 while (!G[u][v0]) {
   L.emplace_back(v, d = X[v]);
   if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
     c = color(u, L[a].first, c);
   else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
     color(u, L[a].first, L[a].second);
   else if (vst[d]) break
   else vst[d] = 1, v = C[u][d];
 if (!G[u][v0]) {
   for (; v; v = flip(v, c, d), swap(c, d));
   if (C[u][c0]) { a = int(L.size()) - 1;
    while (--a >= 0 \&\& L[a].second != c)
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
   } else t--;
```

4 Matching & Flow

4.1 Kuhn Munkres

```
class KM {
private:
static constexpr 1ld INF = 1LL << 60;</pre>
vector<lld> hl,hr,slk;
vector<int> fl,fr,pre,qu;
vector<vector<lld>> w;
vector<bool> v1,vr;
int n, ql, qr;
bool check(int x) {
 if (vl[x] = true, fl[x] != -1)
  return vr[qu[qr++] = fl[x]] = true;
 while (x != -1) swap(x, fr[fl[x] = pre[x]]);
 return false;
 void bfs(int s) {
 fill(slk.begin(), slk.end(), INF);
 fill(v1.begin(), v1.end(), false);
 fill(vr.begin(), vr.end(), false);
  ql = qr = 0;
  qu[qr++] = s;
 vr[s] = true;
  while (true) {
  11d d;
   while (ql < qr) {</pre>
```

```
if(!vl[x]\&slk[x]>=(d=hl[x]+hr[y]-w[x][y])){
      if (pre[x] = y, d) slk[x] = d;
      else if (!check(x)) return;
    }
   d = INF;
   for (int x = 0; x < n; ++x)
    if (!v1[x] \&\& d > s1k[x]) d = s1k[x];
   for (int x = 0; x < n; ++x) {
    if (v1[x]) h1[x] += d;
    else slk[x] -= d;
    if (vr[x]) hr[x] -= d;
   for (int x = 0; x < n; ++x)
    if (!v1[x] && !slk[x] && !check(x)) return;
 }
public:
 void init( int n_ ) {
  n = n_{;} qu.resize(n);
  fl.clear(); fl.resize(n, -1);
fr.clear(); fr.resize(n, -1);
hr.clear(); hr.resize(n); hl.resize(n);
  w.clear(); w.resize(n, vector<lld>(n));
  slk.resize(n); pre.resize(n);
  vl.resize(n); vr.resize(n);
 void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 1ld solve() {
  for (int i = 0; i < n; ++i)</pre>
   hl[i] = *max_element(w[i].begin(), w[i].end());
  for (int i = 0; i < n; ++i) bfs(i);</pre>
  11d res = 0;
  for (int i = 0; i < n; ++i) res += w[i][f1[i]];</pre>
  return res;
 }
} km;
4.2 Bipartite Matching
class BipartiteMatching{
private:
 vector<int> X[N], Y[N];
 int fX[N], fY[N], n;
 bitset<N> walked;
 bool dfs(int x)
  for(auto i:X[x]){
   if(walked[i])continue;
   walked[i]=1
   if(fY[i]==-1||dfs(fY[i])){
    fY[i]=x;fX[x]=i;
    return 1:
  return 0;
public:
 void init(int _n){
  n=_n; walked.reset();
  for(int i=0;i<n;i++){</pre>
   X[i].clear();Y[i].clear();
   fX[i]=fY[i]=-1;
  }
 void add_edge(int x, int y){
  X[x].push_back(y); Y[y].push_back(y);
 int solve(){
  int cnt = 0;
  for(int i=0;i<n;i++){</pre>
   walked.reset();
   if(dfs(i)) cnt++;
  // return how many pair matched
  return cnt;
};
4.3 General Graph Matching
```

namespace matching {

int fa[kN], pre[kN], match[kN], s[kN], v[kN];

for (int x = 0, y = qu[q1++]; x < n; ++x)

bool SPFA(int u){

```
vector<int> g[kN];
                                                                 if (onstk[u]) return true;
queue<int> q;
                                                                 stk.PB(u); onstk[u] = 1;
                                                                 for (int v=0; v<n; v++){</pre>
void Init(int n) {
for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
                                                                  if (u != v && match[u] != v && !onstk[v]){
                                                                   int m = match[v]
                                                                   if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
void AddEdge(int u, int v) {
                                                                    dis[m] = dis[u] - edge[v][m] + edge[u][v];
                                                                    onstk[v] = 1;
g[u].push_back(v);
 g[v].push_back(u);
                                                                     stk.PB(v)
                                                                    if (SPFA(m)) return true;
int Find(int u) {
                                                                    stk.pop_back();
return u == fa[u] ? u : fa[u] = Find(fa[u]);
                                                                    onstk[v] = 0;
int LCA(int x, int y, int n) {
static int tk = 0; tk++;
                                                                 onstk[u] = 0; stk.pop_back();
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
                                                                 return false;
   if (v[x] == tk) return x;
                                                                int solve() { // find a match
   v[x] = tk;
                                                                 for (int i=0; i<n; i+=2){</pre>
   x = Find(pre[match[x]]);
                                                                  match[i] = i+1;
  }
                                                                  match[i+1] = i;
                                                                 while (true){
                                                                  int found = 0;
void Blossom(int x, int y, int 1) {
 while (Find(x) != 1) {
                                                                  for (int i=0; i<n; i++)</pre>
  pre[x] = y, y = match[x];
                                                                   dis[i] = onstk[i] = 0;
 if (s[y] == 1) q.push(y), s[y] = 0;
if (fa[x] == x) fa[x] = 1;
if (fa[y] == y) fa[y] = 1;
                                                                  for (int i=0; i<n; i++){</pre>
                                                                   stk.clear()
                                                                   if (!onstk[i] && SPFA(i)){
  x = pre[y];
                                                                    found = 1;
 }
                                                                    while (SZ(stk)>=2){
                                                                     int u = stk.back(); stk.pop_back();
bool Bfs(int r, int n) {
                                                                     int v = stk.back(); stk.pop_back();
for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;
                                                                     match[u] = v;
                                                                     match[v] = u;
 while (!q.empty()) q.pop();
 q.push(r);
                                                                   }
 s[r] = 0;
 while (!q.empty()) {
  int x = q.front(); q.pop();
                                                                  if (!found) break;
  for (int u : g[x]) {
  if (s[u] == -1) {
                                                                 int ret = 0;
    pre[u] = x, s[u] = 1;
                                                                 for (int i=0; i<n; i++)</pre>
    if (match[u] == n) {
                                                                  ret += edge[i][match[i]];
     for (int a = u, b = x, last; b != n; a = last, b =
                                                                 return ret>>1:
     pre[a])
      last = match[b], match[b] = a, match[a] = b;
                                                               } graph;
     return true;
                                                               4.5 Minimum Cost Circulation
                                                               struct Edge { int to, cap, rev, cost; };
    q.push(match[u]);
    s[match[u]] = 0;
                                                               vector<Edge> g[kN];
   } else if (!s[u] && Find(u) != Find(x)) {
                                                               int dist[kN],
                                                                             pv[kN], ed[kN];
    int 1 = LCA(u, x, n);
                                                               bool mark[kN];
    Blossom(x, u, 1);
                                                               int NegativeCycle(int n) {
                                                                memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
    Blossom(u, x, 1);
  }
                                                                int upd = -1;
                                                                for (int i = 0; i <= n; ++i) {
 return false;
                                                                 for (int j = 0; j < n; ++j) {
                                                                  int idx = 0;
                                                                  for (auto &e : g[j]) {
int Solve(int n) {
                                                                   if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
 int res = 0;
 for (int x = 0; x < n; ++x) {
                                                                    dist[e.to] = dist[j] + e.cost;
  if (match[x] == n) res += Bfs(x, n);
                                                                    pv[e.to] = j, ed[e.to] = idx;
                                                                     if (i == n) {
 return res;
                                                                     upd = j;
}}
                                                                     while(!mark[upd])mark[upd]=1,upd=pv[upd];
                                                                     return upd;
      Minimum Weight Matching (Clique version)
                                                                    }
struct Graph {
 // 0-base (Perfect Match)
                                                                   idx++;
 int n, edge[MXN][MXN];
 int match[MXN], dis[MXN], onstk[MXN];
                                                                 }
 vector<int> stk;
 void init(int _n) {
                                                                return -1;
  for (int i=0; i<n; i++) for (int j=0; j<n; j++)</pre>
                                                               int Solve(int n) {
                                                                int rt = -1, ans = 0;
   edge[i][j] = 0;
                                                                while ((rt = NegativeCycle(n)) >= 0) {
                                                                 memset(mark, false, sizeof(mark));
 void set_edge(int u, int v, int w) {
  edge[u][v] = edge[v][u] = w; 
                                                                 vector<pair<int, int>> cyc;
```

while (!mark[rt]) {

```
cyc.emplace_back(pv[rt], ed[rt]);
  mark[rt] = true;
  rt = pv[rt];
 reverse(cyc.begin(), cyc.end());
 int cap = kInf;
 for (auto &i : cyc) {
  auto &e = g[i.first][i.second];
  cap = min(cap, e.cap);
 for (auto &i : cyc) {
  auto &e = g[i.first][i.second];
  e.cap -= cap;
  g[e.to][e.rev].cap += cap;
  ans += e.cost * cap;
return ans;
```

4.6 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}$.

 - 2. For each edge (x,y,l,u), connect $x\to y$ with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect t o s with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the
 - maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect t o s with capacity ∞ and let the flow from S to T be f'. If $f+f'\neq \sum_{v\in V, in(v)>0}in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching ${\cal M}$ on bipartite graph(X,Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise. 2. DFS from unmatched vertices in X.

 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect $x \to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y \to x$ with (cost,cap)=(-c,1)
 - 3. For each edge with c < 0, sum these cost as K, then increase d(y)
 - by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v) > 0, connect $S \to v$ with (cost, cap) =(0, d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =(0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph

 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$ 2. Construct a max flow model, let ${\cal K}$ be the sum of all weights
 - Connect source $s \to v$, $v \in G$ with capacity K
 - For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity
 - 5. For $v \in G$, connect it with sink $v \to t$ with capacity K + 2T - $(\sum_{e \in E(v)} w(e)) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K |V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge
 - (v,t) with capacity $-p_v.$ 2. Create edge (u,v) with capacity w with w being the cost of choosing \boldsymbol{u} without choosing \boldsymbol{v}
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with capacity c_y .

 2. Create edge (x,y) with capacity c_{xy} .

 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4.7 Dinic

```
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
    int to, rev;
    Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
      int u = bfs.front(); bfs.pop();
       for (auto e: G[u]) {
         if (e.cap <= 0 or lv[e.to]!=-1) continue;
         bfs.push(e.to); lv[e.to] = lv[u] + 1;
      }
    }
    return lv[ed] != -1;
  Cap DFS(int u, Cap f){
  if (u == ed) return f;
    Cap ret = 0;
    for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
       auto &e = G[u][i];
       if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
      Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
      G[e.to][e.rev].cap += nf;
       if (f == 0) return ret;
    if (ret == 0) lv[u] = -1;
    return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
    G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
    st = st_, ed = ed_; Cap ret = 0;
    while (BFS()) {
       idx.assign(n, 0);
      Cap f = DFS(st, numeric_limits<Cap>::max());
       ret += f;
       if (f == 0) break;
     return ret;
```

Minimum Cost Maximum Flow 4.8

```
class MiniCostMaxiFlow{
 using Cap = int; using Wei = int64_t;
 using PCW = pair<Cap,Wei>
 static constexpr Cap INF_CAP = 1 << 30;</pre>
 static constexpr Wei INF_WEI = 1LL<<60;</pre>
private:
 struct Edge{
  int to, back;
  Cap cap; Wei wei;
  Edge() {}
  Edge(int a,int b, Cap c, Wei d):
   to(a),back(b),cap(c),wei(d) {}
 int ori, edd;
 vector<vector<Edge>> G;
 vector<int> fa, wh;
 vector<bool> inq;
 vector<Wei> dis;
 PCW SPFA(){
  fill(inq.begin(),inq.end(),false);
  fill(dis.begin(), dis.end(), INF_WEI);
  queue<int> qq; qq.push(ori);
  dis[ori] = 0;
  while(not qq.empty()){
```

```
int u=qq.front();qq.pop();
                                                                   if (c == -1 \mid | g[i] > g[c]) c = i;
   ing[u] = false;
                                                                 if (c == -1) break;
   for(int i=0;i<SZ(G[u]);++i){</pre>
                                                                 v[s = t, t = c] = true;
for (int i = 0; i < n; ++i) {
    Edge e=G[u][i];
    int v=e.to; Wei d=e.wei;
                                                                   if (del[i] | v[i]) continue;
    if(e.cap <= 0 | |dis[v] <= dis[u] + d)
                                                                   g[i] += w[c][i];
     continue;
    dis[v] = dis[u] + d;
    fa[v] = u, wh[v] = i;
    if (inq[v]) continue;
                                                                 return make_pair(s, t);
    qq.push(v);
    inq[v] = true;
                                                                int mincut(int n) {
                                                                int cut = 1e9;
  }
                                                                 memset(del, false, sizeof(del));
  if(dis[edd]==INF_WEI) return {-1, -1};
                                                                 for (int i = 0; i < n - 1; ++i) {
                                                                 int s, t; tie(s, t) = phase(n);
del[t] = true; cut = min(cut, g[t]);
for (int j = 0; j < n; ++j) {</pre>
 Cap mw=INF_CAP;
  for(int i=edd;i!=ori;i=fa[i])
  mw=min(mw,G[fa[i]][wh[i]].cap);
                                                                   w[s][j] += w[t][j]; w[j][s] += w[j][t];
 for (int i=edd;i!=ori;i=fa[i]){
   auto &eg=G[fa[i]][wh[i]];
   eg.cap -= mw;
                                                                 }
   G[eg.to][eg.back].cap+=mw;
                                                                 return cut;
  return {mw, dis[edd]};
                                                                4.11 Dijkstra Cost Flow
public:
                                                               // kN = #(vertices)
void init(int n){
                                                               // MCMF.{Init, AddEdge, MincostMaxflow}
  G.clear();G.resize(n);
                                                               // MincostMaxflow(source, sink, flow_limit, &cost)
  fa.resize(n);wh.resize(n);
                                                               // => flow
  inq.resize(n); dis.resize(n);
                                                               using Pii = pair<int, int>;
                                                               constexpr int kInf = 0x3f3f3f3f, kN = 500;
void add_edge(int st, int ed, Cap c, Wei w){
G[st].emplace_back(ed,SZ(G[ed]),c,w);
                                                               struct Edge {
                                                                int to, rev, cost, flow;
 G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
                                                               }:
                                                               struct MCMF { // 0-based
                                                                int n{}, m{}, s{}, t{};
vector<Edge> graph[kN];
PCW solve(int a, int b){
 ori = a, edd = b;
 Cap cc=0; Wei ww=0;
                                                                 // Larger range for relabeling
 while(true){
                                                                 int64_t dis[kN] = {}, h[kN] = {};
  PCW ret=SPFA();
                                                                 int p[kN] = {};
  if(ret.first==-1) break;
                                                                 void Init(int nn) {
   cc+=ret.first;
                                                                 n = nn;
  ww+=ret.first * ret.second;
                                                                 for (int i = 0; i < n; i++) graph[i].clear();</pre>
                                                                 void AddEdge(int u, int v, int f, int c) {
 return {cc,ww};
                                                                 graph[u].push_back({v
} mcmf;
                                                                   static_cast<int>(graph[v].size()), c, f});
                                                                  graph[v].push_back(
4.9 GomoryHu Tree
                                                                   {u, static_cast<int>(graph[u].size()) - 1,
int g[maxn];
                                                                    -c, 0});
vector<edge> GomoryHu(int n){
 vector<edge> rt;
                                                                 bool Dijkstra(int &max_flow, int64_t &cost) {
for(int i=1;i<=n;++i)g[i]=1;</pre>
                                                                 priority_queue<Pii, vector<Pii>, greater<>> pq;
for(int i=2;i<=n;++i){</pre>
                                                                  fill_n(dis, n, kInf);
  int t=g[i];
                                                                  dis[s] = 0;
 flow.reset(); // clear flows on all edge
                                                                  pq.emplace(0, s);
  rt.push_back({i,t,flow(i,t)});
                                                                  while (!pq.empty()) {
 flow.walk(i); // bfs points that connected to i (use
  edges not fully flow)
                                                                   auto u = pq.top();
                                                                   pq.pop();
 for(int j=i+1;j<=n;++j){</pre>
                                                                   int v = u.second;
  if(g[j]==t && flow.connect(j))g[j]=i; // check if i
                                                                   if (dis[v] < u.first) continue;</pre>
                                                                   for (auto &e : graph[v]) {
    can reach i
                                                                    auto new_dis =
                                                                     dis[v] + e.cost + h[v] - h[e.to];
return rt;
                                                                    if (e.flow > 0 && dis[e.to] > new_dis) {
                                                                     dis[e.to] = new_dis;
                                                                     p[e.to] = e.rev;
4.10 Global Min-Cut
                                                                     pq.emplace(dis[e.to], e.to);
const int maxn = 500 + 5;
                                                                    }
int w[maxn][maxn], g[maxn];
                                                                   }
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
                                                                  if (dis[t] == kInf) return false;
w[x][y] += c; w[y][x] += c;
                                                                  for (int i = 0; i < n; i++) h[i] += dis[i];
                                                                  int d = max_flow;
                                                                  for (int u = t; u != s;
pair<int, int> phase(int n) {
memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
int s = -1, t = -1;
                                                                     u = graph[u][p[u]].to) {
                                                                   auto &e = graph[u][p[u]];
                                                                   d = min(d, graph[e.to][e.rev].flow);
while (true) {
 int c = -1;
                                                                 max_flow -= d;
  for (int i = 0; i < n; ++i) {
                                                                  cost += int64_t(d) * h[t];
  if (del[i] || v[i]) continue;
                                                                  for (int u = t; u != s;
```

```
u = graph[u][p[u]].to) {
  auto &e = graph[u][p[u]];
  e.flow += d;
  graph[e.to][e.rev].flow -= d;
}
return true;
}
int MincostMaxflow(
  int ss, int tt, int max_flow, int64_t &cost) {
  this->s = ss, this->t = tt;
  cost = 0;
  fill_n(h, n, 0);
  auto orig_max_flow = max_flow;
  while (Dijkstra(max_flow, cost) && max_flow) {}
  return orig_max_flow - max_flow;
}
};
```

5 Math

5.1 $\lfloor \frac{n}{i} \rfloor$ Enumeration

```
T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor
```

5.2 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b==(x/y)*a;
}
```

5.3 Pollard Rho

```
// does not work when n is prime
// return any non-trivial factor
llu pollard_rho(llu n){
    static auto f=[](llu x,llu k,llu m){
        return add(k,mul(x,x,m),m);
    };
    if (!(n&1)) return 2;
    mt19937 rnd(120821011);
    while(true){
        llu y=2,yy=y,x=rnd()%n,t=1;
        for(llu sz=2;t==1;sz<<=1) {
            for(llu i=0;i<sz;++i){
                if(t!=1)break;
                yy=f(yy,x,n);
                t=gcd(yy>y?yy-y:y-yy,n);
        }
        y=yy;
    }
    if(t!=1&&t!=n) return t;
}
```

5.4 Pi Count (Linear Sieve)

```
static constexpr int N = 1000000 + 5;
11d pi[N];
vector<int> primes;
bool sieved[N];
11d cube_root(11d x){
lld s=cbrt(x-static_cast<long double>(0.1));
while(s*s*s <= x) ++s;
return s-1;
11d square_root(11d x){
lld s=sqrt(x-static_cast<long double>(0.1));
while(s*s <= x) ++s;
return s-1;
void init(){
primes.reserve(N)
primes.push_back(1);
for(int i=2;i<N;i++) {</pre>
 if(!sieved[i]) primes.push_back(i);
  pi[i] = !sieved[i] + pi[i-1];
  for(int p: primes) if(p > 1) {
  if(p * i >= N) break;
   sieved[p * i] = true;
   if(p % i == 0) break;
```

```
11d phi(11d m, 11d n) {
 static constexpr int MM = 80000, NN = 500;
 static lld val[MM][NN];
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
 1ld ret = phi(m,n-1)-phi(m/primes[n],n-1);
 if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
 return ret;
11d pi_count(11d);
11d P2(11d m, 11d n) {
1ld sm = square_root(m), ret = 0;
for(lld i = n+1;primes[i]<=sm;i++)</pre>
  ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
 return ret;
11d pi_count(11d m) {
 if(m < N) return pi[m];</pre>
 11d n = pi_count(cube_root(m));
 return phi(m, n) + n - 1 - P2(m, n);
```

5.5 Strling Number

5.5.1 First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1) \dots (x+n-1) = \sum_{k=0}^n S_1(n,k) x^k$$

$$g(x) = x(x+1) \dots (x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.5.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

5.6 Range Sieve

```
const int MAX_SQRT_B = 50000;
const int MAX_L = 200000 + 5;
bool is_prime_small[MAX_SQRT_B], is_prime[MAX_L];
void sieve(lld 1, lld r){ // [1, r)
for(lld i=2;i*i<r;i++) is_prime_small[i] = true;
for(lld i=1;i<r;i++) is_prime[i-1] = true;
if(l==1) is_prime[0] = false;
for(lld i=2;i*i<r;i++){
   if(!is_prime_small[i]) continue;
   for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;
   for(lld j=std::max(2LL, (l+i-1)/i)*i;j<r;j+=i)
        is_prime[j-1]=false;
}
</pre>
```

5.7 Miller Rabin

```
VL convolution(const VI &a, const VI &b) {
 if(x<2)return 0;</pre>
                                                                  // Should be able to handle N <= 10^5, C <= 10^4
 if(!(x&1))return x==2;
                                                                  int sz = 1;
 llu x1=x-1; int t=0;
                                                                  while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 while(!(x1&1))x1>>=1,t++;
                                                                  vector<cplx> v(sz);
 for(llu m:magic)if(witn(m,x1,x,t))return 0;
                                                                  for (int i = 0; i < sz; ++i) {
                                                                  double re = i < a.size() ? a[i] : 0;</pre>
 return 1;
                                                                  double im = i < b.size() ? b[i] : 0;</pre>
                                                                   v[i] = cplx(re, im);
5.8 Extended Euler
    a^b \equiv \begin{cases} a^b \mod \varphi(m) + \varphi(m) & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases}
                                                                  fft(v, sz);
                                                  (\text{mod } m)
                                                                  for (int i = 0; i <= sz / 2; ++i) {
                                                                  int j = (sz - i) & (sz - 1);
                                                                   cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
5.9 Gauss Elimination
                                                                     * cplx(0, -0.25);
void gauss(vector<vector<double>> &d) {
                                                                   if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
 int n = d.size(), m = d[0].size();
                                                                     ].conj()) * cplx(0, -0.25);
 for (int i = 0; i < m; ++i) {
                                                                   v[i] = x;
  int p = -1;
  for (int j = i; j < n; ++j) {</pre>
                                                                  ifft(v, sz);
   if (fabs(d[j][i]) < eps) continue;</pre>
                                                                  VL c(sz);
   if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
                                                                  for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);
  if (p == -1) continue;
  for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]); for (int j = 0; j < n; ++j) {
                                                                VI convolution_mod(const VI &a, const VI &b, int p) {
                                                                 int sz = 1;
  if (i == j) continue;
                                                                  while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
   double z = d[j][i] / d[i][i];
                                                                  vector<cplx> fa(sz), fb(sz);
   for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
                                                                  for (int i = 0; i < (int)a.size(); ++i)</pre>
                                                                   fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 }
                                                                  for (int i = 0; i < (int)b.size(); ++i)</pre>
                                                                   fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
                                                                  fft(fa, sz), fft(fb, sz);
5.10
      Fast Fourier Transform
                                                                  double r = 0.25 / sz;
                                                                  cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
const int mod = 1000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353
                                                                   int j = (sz - i) & (sz - 1);
                                                                   cplx a1 = (fa[i] + fa[j].conj());
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
                                                                   cplx a2 = (fa[i] - fa[j].conj()) * r2;
  static_assert (M1 <= M2 && M2 <= M3);
                                                                   cplx b1 = (fb[i] + fb[j].conj()) * r3;
 constexpr int64_t r12 = modpow(M1, M2-2, M2);
                                                                   cplx b2 = (fb[i] - fb[j].conj()) * r4;
                                                                  if (i != j) {
  cplx c1 = (fa[j] + fa[i].conj());
 constexpr int64_t r13 = modpow(M1, M3-2, M3);
  constexpr int64_t r23 = modpow(M2, M3-2, M3)
 constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
                                                                    cplx c2 = (fa[j] - fa[i].conj()) * r2;
 B = (B - A + M2) * r12 % M2;

C = (C - A + M3) * r13 % M3;
                                                                    cplx d1 = (fb[j] + fb[i].conj()) * r3;
                                                                    cplx d2 = (fb[j] - fb[i].conj()) * r4;
 C = (C - B + M3) * r23 % M3;
                                                                    fa[i] = c1 * d1 + c2 * d2 * r5;
  return (A + B * M1 + C * M1M2) % mod;
                                                                    fb[i] = c1 * d2 + c2 * d1:
                                                                  fa[j] = a1 * b1 + a2 * b2 * r5;
                                                                  fb[j] = a1 * b2 + a2 * b1;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
                                                                  fft(fa, sz), fft(fb, sz);
const double pi = acos(-1);
                                                                  vector<int> res(sz);
                                                                  for (int i = 0; i < sz; ++i) {
cplx omega[maxn + 1];
                                                                  long long a = round(fa[i].re), b = round(fb[i].re),
void prefft() -
                                                                        c = round(fa[i].im);
 for (int i = 0; i <= maxn; i++)</pre>
                                                                   res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
 omega[i] = cplx(cos(2 * pi * j / maxn),
                                                                 }
     sin(2 * pi * j / maxn));
                                                                  return res;
void fft(vector<cplx> &v, int n) {
                                                                }}
 int z = __builtin_ctz(n) - 1;
                                                                 5.11 Chinese Remainder
 for (int i = 0; i < n; ++i) {
  int x = 0, j = 0;
                                                                lld crt(lld ans[], lld pri[], int n){
  for (;(1 << j) < n;++j) x^{=(i >> j & 1)<<(z - j);
                                                                  lld M = 1, ret = 0;
  if (x > i) swap(v[x], v[i]);
                                                                  for(int i=0;i<n;i++) M *= pri[i];</pre>
                                                                  for(int i=0;i<n;i++){</pre>
                                                                  1ld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
 for (int s = 2; s <= n; s <<= 1) {
  int z = s >> 1;
                                                                  ret += (ans[i]*(M/pri[i])%M * iv)%M;
  for (int i = 0; i < n; i += s) {
                                                                  ret %= M;
   for (int k = 0; k < z; ++k) {
  cplx x = v[i + z + k] * omega[maxn / s * k];
                                                                  return ret;
    v[i + z + k] = v[i + k] - x;
    v[i + k] = v[i + k] + x;
                                                                /*
                                                                Another:
                                                                x = a1 \% m1
                                                                x = a2 \% m2
                                                                g = gcd(m1, m2)
void ifft(vector<cplx> &v, int n) {
                                                                assert((a1-a2)%g==0)
fft(v, n); reverse(v.begin() + 1, v.end());
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
                                                                [p, q] = exgcd(m2/g, m1/g)
                                                                return a2+m2*(p*(a1-a2)/g)
```

0 <= x < lcm(m1, m2)

```
*/
                                                                  const LL *data() const { return coef.data(); }
                                                                  LL &operator[](size_t i) { return coef[i]; }
 5.12 Berlekamp Massey
                                                                  const LL &operator[](size_t i)const{return coef[i];}
                                                                  Poly(initializer_list<LL> a) : coef(a) { }
 // x: 1-base, p[]: 0-base
                                                                  explicit Poly(int _n = 1) : coef(_n) { }
 template<size_t N>
                                                                 Poly(const LL *arr, int _n) : coef(arr, arr + _n) {}
Poly(const Poly &p, int _n) : coef(_n) {
 vector<llf> BM(llf x[N], size_t n){
 size_t f[N]={0},t=0;11f d[N];
                                                                  copy_n(p.data(), min(p.n(), _n), data());
 vector<llf> p[N];
 for(size_t i=1,b=0;i<=n;++i) {</pre>
                                                                  Poly& irev(){return reverse(data(),data()+n()),*this;}
  for(size_t j=0;j<p[t].size();++j)</pre>
                                                                  Poly& isz(int _n) { return coef.resize(_n), *this; }
    d[i]+=x[i-j-1]*p[t][j];
                                                                  Poly& iadd(const Poly &rhs) { // n() == rhs.n()
   if(abs(d[i]-=x[i])<=EPS)continue;</pre>
                                                                  fi(0, n()) if ((coef[i]+=rhs[i]) >= P)coef[i]-=P;
   f[t]=i;if(!t){p[++t].resize(i);continue;}
                                                                   return *this;
   vector<llf> cur(i-f[b]-1);
   11f k=-d[i]/d[f[b]];cur.PB(-k);
                                                                  Poly& imul(LL k) {
  for(size_t j=0;j<p[b].size();j++)</pre>
                                                                  fi(0, n()) coef[i] = coef[i] * k % P;
   cur.PB(p[b][j]*k);
                                                                   return *this;
   if(cur.size()<p[t].size())cur.resize(p[t].size());</pre>
  for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];</pre>
                                                                  Poly Mul(const Poly &rhs) const {
  if(i-f[b]+p[b].size()>=p[t].size()) b=t;
                                                                  const int _n = n2k(n() + rhs.n() - 1);
  p[++t]=cur;
                                                                  Poly X(*this, _n), Y(rhs, _n);
ntt(X.data(), _n), ntt(Y.data(), _n);
fi(0, _n) X[i] = X[i] * Y[i] % P;
 return p[t];
                                                                   ntt(X.data(), _n, true);
                                                                   return X.isz(n() + rhs.n() - 1);
 5.13 NTT
 template <int mod, int G, int maxn>
                                                                  Poly Inv() const { // coef[0] != 0
 struct NTT {
                                                                  if (n() == 1) return {ntt.minv(coef[0])};
const int _n = n2k(n() * 2);
Poly Xi = Poly(*this, (n() + 1)/2).Inv().isz(_n);
 static_assert (maxn == (maxn & -maxn));
  int roots[maxn];
 NTT () {
                                                                   Poly Y(*this, _n);
  int r = modpow(G, (mod - 1) / maxn);
                                                                   ntt(Xi.data(), _n), ntt(Y.data(), _n);
  for (int i = maxn >> 1; i; i >>= 1) {
                                                                   fi(0, _n) {
   roots[i] = 1;
                                                                    Xi[i] *= (2 - Xi[i] * Y[i]) % P;
    for (int j = 1; j < i; j++)
                                                                    if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
     roots[i + j] = modmul(roots[i + j - 1], r);
    r = modmul(r, r);
                                                                   ntt(Xi.data(), _n, true);
                                                                   return Xi.isz(n());
  // n must be 2^k, and 0 \ll F[i] \ll mod
                                                                  Poly Sqrt() const { // Jacobi(coef[0], P) = 1
 void inplace_ntt(int n, int F[], bool inv = false) {
                                                                   if (n()==1) return {QuadraticResidue(coef[0], P)};
  for (int i = 0, j = 0; i < n; i++) {
  if (i < j) swap(F[i], F[j]);</pre>
                                                                   Poly X = Poly(*this, (n()+1) / 2).Sqrt().isz(n());
                                                                   return X.iadd(Mul(X.Inv()).isz(n())).imul(P/2+1);
    for (int k = n > 1; (j^* = k) < k; k > = 1);
                                                                  pair<Poly, Poly> DivMod(const Poly &rhs) const {
   for (int s = 1; s < n; s *= 2) {
                                                                  // (rhs.)back() != 0
   for (int i = 0; i < n; i += s * 2) {
                                                                   if (n() < rhs.n()) return {{0}, *this};</pre>
     for (int j = 0; j < s; j++) {
                                                                   const int _n = n() - rhs.n() + 1;
      int a = F[i+j];
                                                                   Poly X(rhs); X.irev().isz(_n);
      int b = modmul(F[i+j+s], roots[s+j]);
                                                                   Poly Y(*this); Y.irev().isz(_n);
      F[i+j] = modadd(a, b); // a + b
                                                                   Poly Q = Y.Mul(X.Inv()).isz(_n).irev();
      F[i+j+s] = modsub(a, b); // a - b
                                                                   X = rhs.Mul(Q), Y = *this
                                                                   fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
                                                                   return {Q, Y.isz(max(1, rhs.n() - 1))};
   if (inv) {
                                                                  Poly Dx() const
   int invn = modinv(n);
                                                                  Poly ret(n() - 1);
   for (int i = 0; i < n; i++)
F[i] = modmul(F[i], invn);</pre>
                                                                   fi(0, ret.n()) ret[i] = (i + 1) * coef[i + 1] % P;
                                                                   return ret.isz(max(1, ret.n()));
    reverse(F + 1, F + n);
                                                                  Poly Sx() const {
                                                                  Poly ret(n() + 1);
                                                                   fi(0, n()) ret[i + 1]=ntt.minv(i + 1)*coef[i] % P;
 const int P=2013265921, root=31;
                                                                   return ret;
 const int MAXN=1<<20;</pre>
NTT<P, root, MAXN> ntt;
                                                                  Poly _tmul(int nn, const Poly &rhs) const {
                                                                  Poly Y = Mul(rhs).isz(n() + nn - 1);
 5.14 Polynomial Operations
                                                                   return Poly(Y.data() + n() - 1, nn);
 using VL = vector<LL>;
#define fi(s, n) for (int i=int(s); i<int(n); ++i)
#define Fi(s, n) for (int i=int(n); i>int(s); --i)
                                                                  VL _eval(const VL &x, const auto up)const{
                                                                  const int _n = (int)x.size();
 int n2k(int n) {
                                                                   if (!_n) return {};
 int sz = 1; while (sz < n) sz <<= 1;</pre>
                                                                   vector<Poly> down(_n * 2);
                                                                   down[1] = DivMod(up[1]).second;
 return sz;
                                                                   fi(2,_n*2) down[i]=down[i/2].DivMod(up[i]).second;
 template<int MAXN, LL P, LL RT> // MAXN = 2^k
                                                                   /* down[1] = Poly(up[1]).irev().isz(n()).Inv().irev()
                                                                       tmul(_n, *this)
 struct Poly { // coefficients in [0, P)
                                                                   fi(2, _n * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
 static NTT<MAXN, P, RT> ntt;
                                                                     1, down[i / 2]); */
 VL coef;
 int n() const { return coef.size(); } // n()>=1
                                                                  VL y(_n);
                                                                   fi(0, _n) y[i] = down[_n + i][0];
 LL *data() { return coef.data(); }
```

```
*z = (x0y0 + x1y1 , x0y1 + x1y0 )
  return y:
                                                             * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
static vector<Poly> _tree1(const VL &x) {
                                                             * z' = ((x\theta+x1)(y\theta+y1)), (x\theta-x1)(y\theta-y1))
* z = (1/2) * z''
 const int _n = (int)x.size();
  vector<Poly> up(_n * 2);
                                                             * or convolution:
  fi(0, _n) up[_n + i] = \{(x[i] ? P - x[i] : 0), 1\};
 Fi(0, _n-1) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
                                                             * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                             * and convolution:
  return up:
                                                             * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
VL Eval(const VL&x)const{return _eval(x,_tree1(x));}
                                                            const LL MOD = 1e9+7;
static Poly Interpolate(const VL &x, const VL &y) {
                                                            inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
 const int _n = (int)x.size();
                                                             for( int d = 1 ; d < N ; d <<= 1 ) {</pre>
  vector < Poly > up = _tree1(x), down(_n * 2);
                                                              int d2 = d << 1;
 VL z = up[1].Dx()._eval(x, up);
                                                              for( int s = 0 ; s < N ; s += d2 )
 fi(0, _n) z[i] = y[i] * ntt.minv(z[i]) % P;
                                                               x[ i ] = ta+tb;
                                                                x[ j ] = ta-tb;
   .iadd(down[i * 2 + 1].Mul(up[i * 2]));
                                                                if( x[ i ] >= MOD ) x[ i ] -= MOD;
  return down[1];
                                                                if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
Poly Ln() const { // coef[0] == 1
  return Dx().Mul(Inv()).Sx().isz(n());
                                                             if( inv )
                                                              for( int i = 0 ; i < N ; i++ ) {</pre>
Poly Exp() const \{ // coef[0] == 0 \}
 if (n() == 1) return {1};
Poly X = Poly(*this, (n() + 1)/2).Exp().isz(n());
Poly Y = X.Ln(); Y[0] = P - 1;
                                                               x[ i ] *= inv( N, MOD );
                                                               x[ i ] %= MOD;
  fi(0, n()) if((Y[i] = coef[i] - Y[i]) < 0)Y[i]+=P;
 return X.Mul(Y).isz(n());
                                                                  DiscreteLog
                                                            5.16
Poly Pow(const string &K) const {
                                                            template<typename Int>
                                                            Int BSGS(Int x, Int y, Int M) {
 int nz = 0:
  while (nz < n() && !coef[nz]) ++nz;</pre>
                                                              // x^? \equiv y (mod M)
                                                              Int t = 1, c = 0, g = 1;
 LL nk = 0, nk2 = 0;
                                                              for (Int M_ = M; M_ > 0; M_ >>= 1)
  for (char c : K) {
  nk = (nk * 10 + c - '0') % P;
                                                                g = g * x % M;
  nk2 = nk2 * 10 + c - '0';
                                                              for (g = gcd(g, M); t % g != 0; ++c) {
   if (nk2 * nz >= n()) return Poly(n());
                                                                if (t == y) return c;
   nk2 %= P - 1;
                                                                t = t * x % M;
  if (!nk && !nk2) return Poly({1}, n());
                                                              if (y % g != 0) return -1;
                                                              t /= g, y /= g, M /= g;
Int h = 0, gs = 1;
 Poly X(data() + nz, n() - nz * nk2);
 LL \times 0 = X[0]
                                                              for (; h * h < M; ++h) gs = gs * x % M;
  return X.imul(ntt.minv(x0)).Ln().imul(nk).Exp()
   .imul(ntt.mpow(x0, nk2)).irev().isz(n()).irev();
                                                              unordered_map<Int, Int> bs;
                                                              for (Int s = 0; s < h; bs[y] = ++s)
Poly InvMod(int L) { // (to evaluate linear recursion)
                                                                y = y * x % M;
 Poly R{1, 0}; // *this * R mod x^L = 1 (*this[0] ==
                                                              for (Int s = 0; s < M; s += h) {
                                                                t = t * gs % M;
  for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                                if (bs.count(t)) return c + s + h - bs[t];
  Poly 0 = R.Mul(Poly(data(), min(2 \ll level, n())));
   Poly Q(2 << level); Q[0] = 1;
                                                              return -1;
  for (int j = (1 << level); j < (2 << level); ++j)
Q[j] = (P - O[j]) % P;</pre>
                                                            5.17
                                                                   FloorSum
   R = R.Mul(Q).isz(4 << level);
                                                            // @param n `n < 2^32`
                                                            // @param m `1 <= m < 2^32`
  return R.isz(L);
                                                            // @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
static LL LinearRecursion(const VL&a,const VL&c,LL n){
                                                            llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
 // a_n = \sum_{j=0}^{n} a_{n-j}
                                                             11u ans = 0:
 const int k = (int)a.size();
                                                             while (true)
  assert((int)c.size() == k + 1);
                                                              if (a >= m) {
 Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
                                                               ans += n * (n - 1) / 2 * (a / m); a %= m;
  fi(1, k + 1) C[k - i] = c[i] ? P - c[i] : 0;
 C[k] = 1
                                                              if (b >= m) {
 while (n) {
                                                               ans += n * (b / m); b %= m;
  if (n % 2) W = W.Mul(M).DivMod(C).second;
   n /= 2, M = M.Mul(M).DivMod(C).second;
                                                              llu y_max = a * n + b;
                                                              if (y_max < m) break;</pre>
 LL ret = 0:
                                                              // y_max < m * (n + 1)
  fi(0, k) ret = (ret + W[i] * a[i]) % P;
                                                              // floor(y_max / m) <= n
 return ret;
                                                              n = (11u)(y_max / m), b = (11u)(y_max % m);
}
                                                              swap(m, a);
#undef fi
                                                             return ans;
#undef Fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
                                                            lld floor_sum(lld n, lld m, lld a, lld b) {
template<> decltype(Poly_t::ntt) Poly_t::ntt = {};
                                                             llu ans = 0;
                                                             if (a < 0) {
5.15 FWT
                                                              11u \ a2 = (a \% m + m) \% m;
/* xor convolution:
                                                              ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
* x = (x0, x1) , y = (y0, y1)
                                                              a = a2;
```

```
if (b < 0) {
 11\dot{u} b2 = (b % m + m) % m;
 ans -= 1ULL * n * ((b2 - b) / m);
b = b2
return ans + floor_sum_unsigned(n, m, a, b);
```

5.18 ExtendedFloorSum

```
g(a,b,c,n) = \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor
                              a \geq c \vee b \geq c
                                                                                                                     n < 0 \lor a = 0
                               \begin{bmatrix} \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), \end{bmatrix} 
                                                                                                                       otherwise
h(a,b,c,n) = \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2
                              \left( \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right)
                                +\lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1)
                                +h(a\bmod c,b\bmod c,c,n)
                                +2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)
                             \left\{+2\left\lfloor\frac{b}{c}\right\rfloor\cdot f(a\ \mathsf{mod}\ c,b\ \mathsf{mod}\ c,c,n),\right.
                                                                                                                      a \geq c \vee b \geq c
```

nm(m+1) - 2g(c, c-b-1, a, m-1)

-2f(c, c-b-1, a, m-1) - f(a, b, c, n), otherwise

5.19 Quadratic residue

```
struct S {
 int MOD, w;
 int64_t x, y;
 S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
 : MOD(m), w(w_), x(x_), y(y_) {}
S operator*(const S &rhs) const {
  int w_{-} = w;
  if (w_ == -1) w_ = rhs.w;
  assert(w_ != -1 and w_ == rhs.w);
  return { MOD, w_,
(x * rhs.x + y * rhs.y % MOD * w) % MOD,
   (x * rhs.y + y * rhs.x) % MOD };
};
int get_root(int n, int P) {
  if (P == 2 or n == 0) return n;
  if (qpow(n, (P - 1) / 2, P) != 1) return -1;
  auto check = [&](int x) {
  return qpow(x, (P - 1) / 2, P); };
if (check(n) == P-1) return -1;
  int64_t a; int w; mt19937 rnd(7122);
  do { a = rnd() % P;
w = ((a * a - n) % P + P) % P;
  } while (check(w) != P - 1);
  return qpow(S(P, w, a, 1), (P + 1) / 2).x;
```

5.20 De-Bruijn

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
  if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
 } else {
 aux[t] = aux[t - p];
db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {
   aux[t] = i;
   db(t + 1, t, n, k);
int de_bruijn(int k, int n) {
// return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
if (k == 1) {
 res[0] = 0;
  return 1;
```

```
for (int i = 0; i < k * n; i++) aux[i] = 0;
db(1, 1, n, k);
return sz;
```

5.21 Simplex Construction

Standard form: maximize $\sum_{1 \leq i \leq n} c_i x_i$ such that for all $1 \leq j \leq m$, $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$ and $x_i \ge 0$ for all $1 \le i \le n$.

- 1. In case of minimization, let $c_i^\prime = -c_i$
- 2. $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j$
- $3. \sum_{1 < i < n} A_{ji} x_i = b_j$
 - $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$
 - $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

5.22 Simplex

 $n < 0 \lor a = 0$

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return VD(n, -inf) if the solution doesn't exist // return VD(n, +inf) if the solution is unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
 double inv = 1.0 / d[r][s];
 for (int i = 0; i < m + 2; ++i)
  for (int j = 0; j < n + 2; ++j)
   if (i != r && j != s)
    d[i][j] -= d[r][j] * d[i][s] * inv;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;
 d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
 int x = m + z;
 while (true) {
  int s = -1;
  for (int i = 0; i <= n; ++i) {
  if (!z && q[i] == -1) continue;
   if (s == -1) \mid d[x][i] < d[x][s]) s = i;
  if (d[x][s] > -eps) return true;
  int r = -1;
  for (int i = 0; i < m; ++i) {</pre>
   if (d[i][s] < eps) continue;
if (r == -1 || \</pre>
     d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
  pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
 m = b.size(), n = c.size();
 d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)</pre>
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)
  p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
 for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i]; q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
 for (int i = 1; i < m; ++i)
if (d[i][n + 1] < d[r][n + 1]) r = i;
 if (d[r][n + 1] < -eps) {</pre>
  pivot(r, n);
  if (!phase(1) \mid | d[m + 1][n + 1] < -eps)
   return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
   int s = min_element(d[i].begin(), d[i].end() - 1)
        - d[i].begin();
```

```
pivot(i, s);
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)
  if (p[i] < n) x[p[i]] = d[i][n + 1];
 return x:
5.23 Charateristic Polynomial
vector<vector<int>> Hessenberg(const vector<vector<int
     >> &A) {
 int N = A.size();
 vector<vector<int>> H = A;
 for (int i = 0; i < N - 2; ++i) {
  if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {
    if (H[j][i]) {
     for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
      for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k
     ][j]);
     break:
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
   * H[i + 1][k] * (kP - coef)) % kP;
for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
  }
 return H;
vector<int> CharacteristicPoly(const vector<vector<int
    >> &A) {
 int N = A.size();
 auto H = Hessenberg(A);
for (int i = 0; i < N; ++i) {</pre>
  for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
 vector<vector<int>>> P(N + 1, vector<int>(N + 1));
 P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {</pre>
  P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j -
    1];
  int val = 1;
  for (int j = i - 1; j >= 0; --j) {
  int coef = 1LL * val * H[j][i - 1] % kP;
   for (int k = 0; k \le j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 if (N & 1) {
  for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
 return P[N];
5.24 Partition Number
int b = sqrt(n);
ans[0] = tmp[0] = 1;
for (int i = 1; i <= b; i++) {
for (int rep = 0; rep < 2; rep++)
for (int j = i; j <= n - i * i; j++)
modadd(tmp[j], tmp[j-i]);
for (int j = i * i; j <= n; j++)</pre>
  modadd(ans[j], tmp[j - i * i]);
```

6 Geometry

6.1 Basic Geometry

```
using coord_t = int;
using Real = double;
using Point = std::complex<coord_t>;
int sgn(coord_t x) {
return (x > 0) - (x < 0); }
coord_t dot(Point a, Point b) {
 return real(conj(a) * b); }
coord_t cross(Point a, Point b) {
 return imag(conj(a) * b); }
int ori(Point a, Point b, Point c) {
 return sgn(cross(b - a, c - a)); }
bool operator<(const Point &a, const Point &b) {</pre>
 return real(a) != real(b)
  ? real(a) < real(b) : imag(a) < imag(b);
int argCmp(Point a, Point b) {
  // -1 / 0 / 1 <-> < / == / > (atan2)
 int qa = (imag(a) == 0
   ? (real(a) < 0 ? 3 : 1) : (imag(a) < 0 ? 0 : 2));
 int qb = (imag(b) == 0
   ? (real(b) < 0 ? 3 : 1) : (imag(b) < 0 ? 0 : 2));
 if (qa != qb) return sgn(qa - qb);
 return sgn(cross(b, a));
template <typename V> Real area(const V & pt) {
 coord_t ret = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 return ret / 2.0;
6.2 2D Convex Hull
template<typename PT>
vector<PT> buildConvexHull(vector<PT> d) {
 sort(ALL(d), [](const PT& a, const PT& b){
   return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
 vector<PT> s(SZ(d)<<1);</pre>
 int o = 0;
 for(auto p: d) {
  while(o>=2 && cross(p-s[o-2], s[o-1]-s[o-2])<=0)
  s[o++] = p;
 for(int i=SZ(d)-2, t = o+1;i>=0;i--){
  while(o>=t&&cross(d[i]-s[o-2],s[o-1]-s[o-2])<=0)
  s[o++] = d[i];
 s.resize(o-1);
 return s;
6.3
     3D Convex Hull
// return the faces with pt indexes
int flag[MXN][MXN];
struct Point{
 ld x,y,z;
 Point operator * (const 1d &b) const {
  return (Point) {x*b,y*b,z*b};}
 Point operator * (const Point &b) const {
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
Point ver(Point a, Point b, Point c) {
  return (b - a) * (c - a);}
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0;
 REP(i,n) REP(j,n) flag[i][j] = 0;
 vector<Face> now
 now.emplace_back(0,1,2);
 now.emplace_back(2,1,0);
for (int i=3; i<n; i++){
  ftop++; vector<Face> next;
  REP(j, SZ(now)) {
   Face& f=now[j]; int ff = 0;
   ld d=(pt[i]-pt[f.a]).dot(
     ver(pt[f.a], pt[f.b], pt[f.c]));
   if (d <= 0) next.push_back(f);</pre>
   if (d > 0) ff=ftop;
   else if (d < 0) ff=-ftop;
flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;</pre>
```

```
REP(j, SZ(now)) {
                                                                   if (dis(v[p], v[i]) < d) {</pre>
   Face& f=now[j]
                                                                    d = dis(v[p], v[i]);
   if (flag[f.a][f.b] > 0 &&
                                                                    found = true;
     flag[f.a][f.b] != flag[f.b][f.a])
    next.emplace_back(f.a,f.b,i);
   if (flag[f.b][f.c] > 0 &&
                                                                 }
     flag[f.b][f.c] != flag[f.c][f.b])
    next.emplace_back(f.b,f.c,i);
                                                               if (found) rebuild_m(i + 1);
   if (flag[f.c][f.a] > 0 &&
                                                                else m[kx][ky][kz] = i;
     flag[f.c][f.a] != flag[f.a][f.c])
    next.emplace_back(f.c,f.a,i);
                                                               return d;
                                                             }
 now=next:
                                                             6.7
                                                                  Simulated Annealing
 return now;
                                                             11f anneal() {
                                                              mt19937 rnd_engine( seed );
     2D Farthest Pair
                                                               uniform_real_distribution< llf > rnd( 0, 1 );
6.4
                                                               const llf dT = 0.001;
// stk is from convex hull
                                                               // Argument p
n = (int)(stk.size());
                                                               1lf S_cur = calc( p ), S_best = S_cur;
for ( 1lf T = 2000 ; T > EPS ; T -= dT ) {
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {</pre>
                                                                // Modify p to p_prime
 while(abs(cross(stk[i+1]-stk[i],
                                                                const llf S_prime = calc( p_prime );
   stk[(pos+1)%n]-stk[i])) >
                                                               const llf delta_c = S_prime - S_cur;
llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
   abs(cross(stk[i+1]-stk[i],
   stk[pos]-stk[i]))) pos = (pos+1)%n;
                                                               if ( rnd( rnd_engine ) <= prob )</pre>
 ans = max({ans, dis(stk[i], stk[pos]),
                                                                S_cur = S_prime, p = p_prime;
  dis(stk[i+1], stk[pos])});
                                                                if ( S_prime < S_best ) // find min</pre>
                                                                 S_best = S_prime, p_best = p_prime;
6.5 2D Closest Pair
                                                               return S_best;
struct cmp_y {
 bool operator()(const P& p, const P& q) const {
  return p.y < q.y;</pre>
                                                              6.8 Half Plane Intersection
                                                             // NOTE: Point is complex<Real>
multiset<P, cmp_y> s;
                                                              // cross(pt-line.st, line.dir)<=0 <-> pt in half plane
void solve(P a[], int n) {
                                                             struct Line {
 sort(a, a + n, [](const P& p, const P& q) {
                                                               Point st, ed;
  return tie(p.x, p.y) < tie(q.x, q.y);</pre>
                                                                Point dir;
                                                               Line (Point _s, Point _e)
 11f d = INF; int pt = 0;
                                                                 : st(_s), ed(_e), dir(_e - _s) {}
 for (int i = 0; i < n; ++i) {
 while (pt < i \text{ and } a[i].x - a[pt].x >= d)
                                                             bool operator<(const Line &lhs, const Line &rhs) {</pre>
   s.erase(s.find(a[pt++]));
                                                                if (int cmp = argCmp(lhs.dir, rhs.dir))
  auto it = s.lower_bound(P(a[i].x, a[i].y - d));
                                                                  return cmp == -1;
  while (it != s.end() and it->y - a[i].y < d)
                                                                return ori(lhs.st, lhs.ed, rhs.st) < 0;</pre>
   d = min(d, dis(*(it++), a[i]));
  s.insert(a[i]);
                                                             Point intersect(const Line &A, const Line &B) {
                                                               Real t = cross(B.st - A.st, B.dir) /
                                                                cross(A.dir, B.dir)
                                                                return A.st + t * A.dir;
6.6 kD Closest Pair (3D ver.)
                                                             Real HPI(vector<Line> &lines) {
11f solve(vector<P> v) {
                                                                sort(lines.begin(), lines.end());
 shuffle(v.begin(), v.end(), mt19937());
                                                                deque<Line> que;
 unordered_map<lld, unordered_map<lld,
                                                                deque<Point> pt;
  unordered_map<lld, int>>> m;
                                                                que.push_back(lines[0]);
 llf d = dis(v[0], v[1]);
                                                                for (int i = 1; i < (int)lines.size(); i++) {</pre>
 auto Idx = [&d] (11f x) -> 11d {
                                                                  if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
  return round(x * 2 / d) + 0.1;
                                                                   continue;
 auto rebuild_m = [&m, &v, &Idx](int k) {
                                                             #define POP(L, R) \
  m.clear():
                                                                  while (pt.size() > 0 \
  for (int i = 0; i < k; ++i)
m[Idx(v[i].x)][Idx(v[i].y)]</pre>
                                                                    && ori(L.st, L.ed, pt.back()) < 0) \
                                                                  pt.pop_back(), que.pop_back(); \
while (pt.size() > 0 \
    [Idx(v[i].z)] = i;
 }; rebuild_m(2);
                                                                    && ori(R.st, R.ed, pt.front()) < 0) \
 for (size_t i = 2; i < v.size(); ++i) {
                                                                    pt.pop_front(), que.pop_front();
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
                                                                  POP(lines[i], lines[i]);
     kz = Idx(v[i].z); bool found = false;
                                                                  pt.push_back(intersect(que.back(), lines[i]));
  for (int dx = -2; dx <= 2; ++dx) {
                                                                  que.push_back(lines[i]);
   const 11d nx = dx + kx
   if (m.find(nx) == m.end()) continue;
                                                               POP(que.front(), que.back())
   auto& mm = m[nx];
                                                                if (que.size() <= 1 ||</pre>
   for (int dy = -2; dy <= 2; ++dy) {
                                                                  argCmp(que.front().dir, que.back().dir) == 0)
    const 11d ny = dy + ky;
                                                                  return 0:
    if (mm.find(ny) == mm.end()) continue;
                                                               pt.push_back(intersect(que.front(), que.back()));
    auto& mmm = mm[ny];
                                                                return area(pt);
    for (int dz = -2; dz <= 2; ++dz) {
     const 11d nz = dz + kz;
     if (mmm.find(nz) == mmm.end()) continue;
                                                                   Minkowski Sum
     const int p = mmm[nz];
```

// sign1 = 1 for outer tang, -1 for inter tang

```
vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
                                                               vector<Line> ret;
hull(A), hull(B);
                                                               double d_{sq} = norm2(c1.0 - c2.0);
 vector<pll> C(1, A[0] + B[0]), s1, s2;
                                                               if( d_sq < eps ) return ret;</pre>
 for(int i = 0; i < SZ(A); ++i)
s1.pb(A[(i + 1) % SZ(A)] - A[i]);
                                                               double d = sqrt( d_sq )
                                                               Pt v = (c2.0 - c1.0) / d;
 for(int i = 0; i < SZ(B); i++)
s2.pb(B[(i + 1) % SZ(B)] - B[i]);
                                                               double c = (c1.R - sign1 * c2.R) / d;
                                                               if( c * c > 1 ) return ret;
 for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
                                                               double h = sqrt( max( 0.0 , 1.0 - c * c ) );
                                                               for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
  Pt n = { v.X * c - sign2 * h * v.Y ,
  if (p2 >= SZ(B)
    || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
                                                                 v.Y * c + sign2 * h * v.X };
   C.pb(C.back() + s1[p1++]);
                                                                Pt p1 = c1.0 + n * c1.R;
                                                                Pt p2 = c2.0 + n * (c2.R * sign1);
   C.pb(C.back() + s2[p2++]);
 return hull(C), C;
                                                                if( fabs( p1.X - p2.X ) < eps and
                                                                  fabs( p1.Y - p2.Y ) < eps )
                                                                 p2 = p1 + perp(c2.0 - c1.0);
6.10 Intersection of line and Circle
                                                                ret.push_back( { p1 , p2 } );
vector<pdd> line_interCircle(const pdd &p1,
                                                               return ret;
    const pdd &p2,const pdd &c,const double r){
                                                              }
 pdd ft=foot(p1,p2,c),vec=p2-p1;
 double dis=abs(c-ft);
                                                                     Minimum Covering Circle
 if(fabs(dis-r)<eps) return vector<pdd>{ft};
                                                              template<typename P>
 if(dis>r) return {};
                                                              Circle getCircum(const P &a, const P &b, const P &c){
 vec=vec*sqrt(r*r-dis*dis)/abs(vec);
                                                               Real a1 = a.x-b.x, b1 = a.y-b.y;
 return vector<pdd>{ft+vec,ft-vec};
                                                               Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
                                                               Real a2 = a.x-c.x, b2 = a.y-c.y;
                                                               Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
6.11 Intersection of Polygon and Circle
                                                               Circle cc;
// Divides into multiple triangle, and sum up
                                                               cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
// test by HDU2892
                                                               cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
const double PI=acos(-1);
                                                               cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
                                                               return cc:
 if(abs(pb)<eps) return 0;</pre>
 double S, h, theta;
                                                              template<typename P>
 double a=abs(pb), b=abs(pa), c=abs(pb-pa);
                                                              Circle MinCircleCover(const vector<P>& pts){
 double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
                                                               random_shuffle(pts.begin(), pts.end());
 double cosC = dot(pa,pb) / a / b, C = acos(cosC);
                                                               Circle c = { pts[0], 0 };
 if(a > r){
                                                               for(int i=0;i<(int)pts.size();i++){</pre>
 S = (C/2)*r*r;
                                                                if (dist(pts[i], c.o) <= c.r) continue;</pre>
 h = a*b*sin(C)/c;
                                                                c = { pts[i], 0 };
for (int j = 0; j < i; j++) {
  if (h < r && B < PI/2)
   S = (acos(h/r)*r*r - h*sqrt(r*r-h*h));
                                                                 if(dist(pts[j], c.o) <= c.r) continue;</pre>
                                                                 c.o = (pts[i] + pts[j]) / 2;
 else if(b > r){
                                                                 c.r = dist(pts[i], c.o)
  theta = PI - B - asin(sin(B)/r*a);
                                                                 for (int k = 0; k < j; k++) {
  S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
                                                                  if (dist(pts[k], c.o) <= c.r) continue;</pre>
                                                                  c = getCircum(pts[i], pts[j], pts[k]);
 else S = .5*sin(C)*a*b;
 return S;
                                                                }
double area_poly_circle(const vector<pdd> poly,
                                                               return c;
  const pdd &0,const double r){
 double S=0;
 for(int i=0;i<SZ(poly);++i)</pre>
                                                                    KDTree (Nearest Point)
                                                              6.15
  S+=_area(poly[i]-0,poly[(i+1)%SZ(poly)]-0,r)
                                                              const int MXN = 100005;
    *ori(0,poly[i],poly[(i+1)%SZ(poly)]);
                                                              struct KDTree {
 return fabs(S);
                                                               struct Node {
}
                                                                int x,y,x1,y1,x2,y2;
                                                                int id,f;
Node *L, *R;
6.12 Intersection of Two Circle
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
                                                               } tree[MXN], *root;
 pdd o1 = a.0, o2 = b.0;
                                                               int n;
 double r1 = a.R, r2 = b.R, d2 = abs2(o1 - o2),
                                                               LL dis2(int x1, int y1, int x2, int y2) {
     d = sqrt(d2)
                                                                LL dx = x1-x2, dy = y1-y2;
 if(d < max(r1, r2) - min(r1, r2) \mid | d > r1 + r2)
                                                                return dx*dx+dy*dy;
  return 0;
 pdd u = (o1 + o2) * 0.5
                                                               static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
  + (o1 - o2) * ((r2 * r2 - r1 * r1) / (2 * d2));
                                                               static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 double A = sqrt((r1 + r2 + d) * (r1 - r2 + d)
                                                               void init(vector<pair<int,int>> ip) {
     * (r1 + r2 - d) * (-r1 + r2 + d));
                                                                n = ip.size();
 pdd v = pdd(o1.Y - o2.Y, -o1.X + o2.X) * A
                                                                for (int i=0; i<n; i++) {</pre>
  / (2 * d2);
                                                                 tree[i].id = i;
 p1 = u + v, p2 = u - v;
                                                                 tree[i].x = ip[i].first;
 return 1;
                                                                 tree[i].y = ip[i].second;
                                                                root = build_tree(0, n-1, 0);
     Tangent line of Two Circle
6.13
vector<Line> go(const Cir& c1,
                                                               Node* build_tree(int L, int R, int d) {
  const Cir& c2, int sign1){
                                                                if (L>R) return nullptr;
                                                                int M = (L+R)/2; tree[M].f = d%2;
```

```
nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
                                                                 int C;
  tree[M].x1 = tree[M].x2 = tree[M].x;
                                                                 Cir c[N]
                                                                 bool g[N][N], overlap[N][N];
  tree[M].y1 = tree[M].y2 = tree[M].y;
                                                                 // Area[i] : area covered by at least i circles
double Area[ N ];
  tree[M].L = build_tree(L, M-1, d+1);
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
                                                                 void init(int _C){ C = _C;}
   tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
                                                                 struct Teve {
                                                                  pdd p; double ang; int add;
   tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
                                                                  Teve() {}
                                                                  Teve(pdd _a, double _b, int _c):p(_a), ang(_b), add(
  tree[M].R = build_tree(M+1, R, d+1);
  if (tree[M].R) {
                                                                  bool operator<(const Teve &a)const
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
                                                                  {return ang < a.ang;}
                                                                 }eve[N * 2];
   tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
                                                                 // strict: x = 0, otherwise x = -1
                                                                 bool disjuct(Cir &a, Cir &b, int x)
   tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
                                                                 {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
                                                                 bool contain(Cir &a, Cir &b, int x)
  return tree+M;
                                                                 \{return sign(a.R - b.R - abs(a.0 - b.0)) > x;\}
 int touch(Node* r, int x, int y, LL d2){
                                                                 bool contain(int i, int j) {
 LL dis = sqrt(d2)+1;
                                                                  /* c[j] is non-strictly in c[i]. */
  if (x<r->x1-dis || x>r->x2+dis ||
                                                                  return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c[i].R) \mid c[i].R - c[i].R)
                                                                     [j].R) == 0 && i < j)) && contain(c[i], c[j], -1);
    y<r->y1-dis || y>r->y2+dis)
   return 0:
                                                                 void solve(){
  return 1:
                                                                  fill_n(Area, C + 2, 0);
                                                                  for(int i = 0; i < C; ++i)
 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
 if (!r || !touch(r, x, y, md2)) return;
LL d2 = dis2(r->x, r->y, x, y);
                                                                   for(int j = 0; j < C; ++j)
                                                                    overlap[i][j] = contain(i, j);
  if (d2 < md2 || (d2 == md2 && mID < r->id)) {
                                                                  for(int i = 0; i < C; ++i)
  mID = r -> id;
                                                                   for(int j = 0; j < C; ++j)
                                                                    g[i][j] = !(overlap[i][j] || overlap[j][i] ||
   md2 = d2;
                                                                       disjuct(c[i], c[j], -1));
  // search order depends on split dim
                                                                  for(int i = 0; i < C; ++i){</pre>
  if ((r->f == 0 \&\& x < r->x) ||
                                                                   int E = 0, cnt = 1;
    (r->f == 1 \&\& y < r->y)) {
                                                                   for(int j = 0; j < C; ++j)</pre>
                                                                    if(j != i && overlap[j][i])
   nearest(r->L, x, y, mID, md2);
   nearest(r->R, x, y, mID, md2);
                                                                     ++cnt;
                                                                   for(int j = 0; j < C; ++j)
if(i != j && g[i][j]) {</pre>
  } else {
   nearest(r->R, x, y, mID, md2);
   nearest(r->L, x, y, mID, md2);
                                                                     pdd aa, bb;
                                                                     CCinter(c[i], c[j], aa, bb);
llf A = atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);
 int query(int x, int y) {
                                                                      11f B = atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);
  int id = 1029384756;
                                                                     eve[E++] = Teve(bb,B,1), eve[E++]=Teve(aa,A,-1);
  LL d2 = 102938475612345678LL;
                                                                     if(B > A) ++cnt;
  nearest(root, x, y, id, d2);
                                                                   if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
  return id;
} tree;
                                                                    sort(eve, eve + E);
                                                                    eve[E] = eve[0];
6.16
      Rotating Sweep Line
                                                                    for(int j = 0; j < E; ++j){}
                                                                     cnt += eve[j].add;
void rotatingSweepLine(pair<int, int> a[], int n) {
 vector<pair<int, int>> 1;
                                                                     Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
 1.reserve(n * (n - 1) / 2)
                                                                      double theta = eve[j + 1].ang - eve[j].ang;
 for (int i = 0; i < n; ++i)
                                                                     if (theta < 0) theta += 2. * pi:</pre>
  for (int j = i + 1; j < n; ++j)
                                                                     Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
   1.emplace_back(i, j);
 sort(l.begin(), l.end(), [&a](auto &u, auto &v){
  1ld udx = a[u.first].first - a[u.second].first;
  lld udy = a[u.first].second - a[u.second].second;
lld vdx = a[v.first].first - a[v.second].first;
  11d vdy = a[v.first].second - a[v.second].second;
  if (udx == 0 \text{ or } vdx == 0) \text{ return not } udx == 0;
                                                                     Stringology
  int s = sgn(udx * vdx);
  return udy * vdx * s < vdy * udx * s;
                                                                7.1 Hash
 });
                                                                class Hash {
 vector<int> idx(n), p(n);
 iota(idx.begin(), idx.end(), 0);
                                                                 private:
                                                                  static constexpr int P = 127, Q = 1051762951;
 sort(idx.begin(), idx.end(), [&a](int i, int j){
 return a[i] < a[j]; });
for (int i = 0; i < n; ++i) p[idx[i]] = i;
                                                                  vector<int> h, p;
                                                                 public:
                                                                  void init(const string &s){
 for (auto [i, j]: 1) {
                                                                   h.assign(s.size()+1, 0); p.resize(s.size()+1);
  // do here
                                                                   for (size_t i = 0; i < s.size(); ++i)</pre>
  swap(p[i], p[j]);
                                                                    h[i + 1] = add(mul(h[i], P), s[i]);
  idx[p[i]] = i, idx[p[j]] = j;
                                                                   generate(p.begin(), p.end(),[x=1,y=1,this]()
                                                                     mutable{y=x;x=mul(x,P);return y;});
6.17 Circle Cover
                                                                  int query(int 1, int r){ // 1-base (1, r]
                                                                   return sub(h[r], mul(h[1], p[r-1]));}
const int N = 1021;
struct CircleCover {
```

7.2 Suffix Arrau namespace sfxarray { bool t[maxn * 2]; int hi[maxn], rev[maxn]; int _s[maxn * 2], sa[maxn * 2], c[maxn * 2]; int x[maxn], p[maxn], q[maxn * 2]; // sa[i]: sa[i]-th suffix is the \ // i-th lexigraphically smallest suffix. // hi[i]: longest common prefix \ // of suffix sa[i] and suffix sa[i - 1]. void pre(int *sa, int *c, int n, int z) { memset(sa, 0, sizeof(int) * n); memcpy(x, c, sizeof(int) * z); void induce(int *sa,int *c,int *s,bool *t,int n,int z){ memcpy(x + 1, c, sizeof(int) * (z - 1)); for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] - 1]) sa[x[s[sa[i] - 1]]++] = sa[i] - 1;memcpy(x, c, sizeof(int) * z); for (int i = n - 1; $i \ge 0$; --i) if (sa[i] && t[sa[i] - 1]) sa[--x[s[sa[i] - 1]]] = sa[i] - 1;void sais(int *s, int *sa, int *p, int *q, bool *t, int *c, int n, int z) { bool uniq = t[n - 1] = true; int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1; memset(c, 0, sizeof(int) * z); for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2; for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i]; if (uniq) { for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i; return: for (int i = n - 2; i >= 0; --i) t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);pre(sa, c, n, z); for (int i = 1; i <= n - 1; ++i)</pre> if (t[i] && !t[i - 1]) sa[--x[s[i]]] = p[q[i] = nn++] = i;induce(sa, c, s, t, n, z); for (int i = 0; i < n; ++i) if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) { bool neq = last < 0 ||</pre> memcmp(s + sa[i], s + last,(p[q[sa[i]] + 1] - sa[i]) * sizeof(int)); ns[q[last = sa[i]]] = nmxz += neq; }} sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1); pre(sa, c, n, z); for (int i = nn - 1; i >= 0; --i) sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];induce(sa, c, s, t, n, z); void build(const string &s) { for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre> _s[(int)s.size()] = 0; // s shouldn't contain 0 sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256); for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];</pre> for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre> int ind = 0; hi[0] = 0; for (int i = 0; i < (int)s.size(); ++i) { if (!rev[i]) {</pre> ind = 0; continue: while (i + ind < (int)s.size() && \</pre> s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;hi[rev[i]] = ind ? ind-- : 0; }} 7.3 Suffix Automaton struct SuffixAutomaton { struct node int ch[K], len, fail, cnt, indeg; $node(int L = 0) : ch{}, len(L), fail(0), cnt(0),$

return f; vector < int > f = kmp(t), r

return res;

```
indeg(0) {}
} st[N];
int root, last, tot;
```

```
void extend(int c) {
  int cur = ++tot;
  st[cur] = node(st[last].len + 1);
  while (last && !st[last].ch[c]) {
    st[last].ch[c] = cur;
    last = st[last].fail;
  if (!last) {
    st[cur].fail = root;
  } else {
    int q = st[last].ch[c];
    if (st[q].len == st[last].len + 1) {
      st[cur].fail = q;
    } else {
      int clone = ++tot;
      st[clone] = st[q];
st[clone].len = st[last].len + 1;
      st[st[cur].fail = st[q].fail = clone].cnt = 0;
      while (last && st[last].ch[c] == q) {
        st[last].ch[c] = clone;
        last = st[last].fail;
      }
  st[last = cur].cnt += 1;
 void init(const char* s) {
  root = last = tot = 1;
  st[root] = node(0);
  for (char c; c = *s; ++s) extend(c - 'a');
 int q[N];
 void dp() +
  for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
  int head = 0, tail = 0;
  for (int i = 1; i <= tot; i++)</pre>
    if (st[i].indeg == 0) q[tail++] = i;
  while (head != tail) {
    int now = q[head++];
    if (int f = st[now].fail) {
      st[f].cnt += st[now].cnt;
      if (--st[f].indeg == 0) q[tail++] = f;
  }
 int run(const char* s) {
  int now = root;
  for (char c; c = *s; ++s)
    if (!st[now].ch[c -= 'a']) return 0;
    now = st[now].ch[c];
  return st[now].cnt;
 }
} SAM;
7.4 KMP
vector<int> kmp(const string &s) {
 vector<int> f(s.size(), 0);
 /* f[i] = length of the longest prefix
   (excluding s[0:i]) such that it coincides with the suffix of s[0:i] of the same length */
 /* i + 1 - f[i] is the length of the
   smallest recurring period of s[0:i] */
 int k = 0;
 for (int i = 1; i < (int)s.size(); ++i) {</pre>
  while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
  if (s[i] == s[k]) ++k;
  f[i] = k;
vector<int> search(const string &s, const string &t) {
 // return 0-indexed occurrence of t in s
 for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
  while(k > 0 && (k==(int)t.size() \mid \mid s[i]!=t[k]))
   k = f[k - 1]
  if (s[i] == t[k]) ++k;
```

if (k == (int)t.size()) r.push_back(i-t.size()+1);

7.5 Z value

| }

```
char s[MAXN];
int len,z[MAXN];
void Z_value() {
  int i,j,left,right;
  z[left=right=0]=len;
  for(i=1;i<len;i++) {
    j=max(min(z[i-left],right-i),0);
    for(;i+j<len&&s[i+j]==s[j];j++);
    if(i+(z[i]=j)>right)right=i+z[left=i];
  }
}
```

7.6 Manacher

```
int z[maxn];
int manacher(const string& s) {
    string t = ".";
    for(char c: s) t += c, t += '.';
    int l = 0, r = 0, ans = 0;
    for (int i = 1; i < t.length(); ++i) {
        z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
        while (i - z[i] >= 0 && i + z[i] < t.length()) {
        if(t[i - z[i]] == t[i + z[i]]) ++z[i];
        else break;
    }
    if (i + z[i] > r) r = i + z[i], l = i;
}
for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
    return ans;
}</pre>
```

7.7 Lexico Smallest Rotation

```
string mcp(string s){
  int n = s.length();
  s += s;
  int i=0, j=1;
  while (i<n && j<n){
    int k = 0;
    while (k < n && s[i+k] == s[j+k]) k++;
    if (s[i+k] <= s[j+k]) j += k+1;
    else i += k+1;
    if (i == j) j++;
}
int ans = i < n ? i : j;
  return s.substr(ans, n);
}</pre>
```

7.8 BWT

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a
 vector<int> v[ SIGMA ];
 void BWT(char* ori, char* res){
// make ori -> ori + ori
  // then build suffix array
 }
 void iBWT(char* ori, char* res){
  for( int i = 0 ; i < SIGMA ; i ++ )</pre>
   v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )</pre>
   v[`ori[i] - BASE ].push_back( i´);
  vector<int> a;
  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
   for( auto j : v[ i ] ){
  a.push_back( j );
  ori[ ptr ++ ] = BASE + i;
  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
   ptr = a[ ptr ];
  res[ len ] = 0;
} bwt;
```

7.9 Palindromic Tree

```
struct palindromic_tree{
 struct node{
  int next[26],f,len;
  int cnt, num, st, ed;
  node( {\color{red} \textbf{int}} \ 1\text{=}\emptyset) : f(\emptyset), len(1), cnt(\emptyset), num(\emptyset) \ \{
   memset(next, 0, sizeof(next)); }
 vector<node> st:
 vector<char> s;
 int last.n:
 void init(){
  st.clear();s.clear();last=1; n=0;
  st.push_back(0);st.push_back(-1);
  st[0].f=1;s.push_back(-1); }
 int getFail(int x){
  while(s[n-st[x].len-1]!=s[n])x=st[x].f;
  return x;}
 void add(int c){
  s.push_back(c-='a'); ++n;
  int cur=getFail(last);
  if(!st[cur].next[c]){
   int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
   st[now].num=st[st[now].f].num+1;
  last=st[cur].next[c];
  ++st[last].cnt;}
 void dpcnt() {
  for (int i=st.size()-1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
 int size(){ return st.size()-2;}
} pt;
int main() {
 string s; cin >> s; pt.init();
 for (int i=0; i<SZ(s); i++) {</pre>
  int prvsz = pt.size(); pt.add(s[i]);
  if (prvsz != pt.size()) {
   int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [1,r]: s.substr(1, r-1+1)
 return 0;
```

8 Misc

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|{\rm det}(\tilde{L}_{rr})|.$

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\dots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \le k \le n$.

8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W\subseteq X, |W|\leq |N_G(W)|$

8.1.7 Euler's planar graph formula

```
V - E + F = C + 1, E \le 3V - 6(?)
```

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Lucas's theorem

```
 \binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p} \text{, where } m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,  and n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0.
```

8.1.10 Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2)$, find maximum $S\in I_1\cap I_2$. For each iteration, build the directed graph and find a shortest path from s to t.

- $s \to x : S \sqcup \{x\} \in I_1$ • $x \to t : S \sqcup \{x\} \in I_2$
- $y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1$ (y is in the unique circuit of $S \sqcup \{x\}$)
- $x \to y : S \setminus \{y\} \sqcup \{x\} \in I_2$ (y is in the unique circuit of $S \sqcup \{x\}$)

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|.$ In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

8.2 DP-opt Condition

8.2.1 totally monotone (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

8.2.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

8.3 Convex 1D/1D DP

```
struct segment {
int i, 1, r;
segment() {}
segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
dp[0] = 0;
deque<segment> dq; dq.push_back(segment(0, 1, n));
for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i);</pre>
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().l = i + 1
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
  f(i, dq.back().1) < f(dq.back().i, dq.back().1))
    dq.pop_back();
  if (dq.size())
   int d = 1 << 20, c = dq.back().1;</pre>
   while (d \gg 1) if (c + d \ll d, back().r)
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
  if (seg.1 <= n) dq.push_back(seg);</pre>
```

8.4 ConvexHull Optimization

```
struct Line {
  mutable int64_t a, b, p;
  bool operator<(const Line &rhs) const { return a < rhs
     .a; }
  bool operator<(int64_t x) const { return p < x; }
};
struct DynamicHull : multiset<Line, less<>> {
```

```
static const int64_t kInf = 1e18;
 bool Isect(iterator x, iterator y)
  auto Div = [](int64_t a, int64_t b) {
    return a / b - ((a ^ b) < 0 && a % b); }
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(int64_t a, int64_t b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
  if (x != begin() \&\& Isect(--x, y)) Isect(x, y = erase)
    (y))
  while ((y = x) != begin() \&\& (--x)->p >= y->p) Isect(
    x, erase(y));
 int64_t Query(int64_t x) {
  auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
};
```

8.5 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

8.6 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
 dfn[u]=low[u]=++dfs_idx;
 for(int i=0;i<(int)init_g[u].size();i++){</pre>
  int v=init_g[u][i];
  if(v==par[u]) continue;
  if(!dfn[v]){
   par[v]=u;
   tarjan(v);
   low[u]=min(low[u],low[v]);
   if(dfn[u]<low[v]){</pre>
    g[u].push_back(v)
    g[v].push_back(u);
  }else{
   low[u]=min(low[u],dfn[v]);
   if(dfn[v]<dfn[u]){
    int temp_v=u;
    bcc_id++
    while(temp_v!=v){
     g[bcc_id+n].push_back(temp_v);
     g[temp_v].push_back(bcc_id+n);
     temp_v=par[temp_v];
    g[bcc_id+n].push_back(v);
    g[v].push_back(bcc_id+n);
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
void dfs(int u,int fa){
 if(u<=n){
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
   dfs(v,u);
   memset(tp,0x8f,sizeof tp);
   if(v<=n){
```

tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);

tp[1]=max(

```
dp[u][0]+dp[v][0]+1
                                                                    ret.push_back( 2 );
     dp[u][1]+max(dp[v][0],dp[v][1])
   }else{
    tp[0]=dp[u][0]+dp[v][0];
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
   dp[u][0]=tp[0],dp[u][1]=tp[1];
 }else{
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
                                                                 }
   if(v==fa) continue;
   dfs(v,u);
  min_dp[0][0]=0;
min_dp[1][1]=1;
min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f3;
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
   memset(tmp,0x8f,sizeof tmp);
   tmp[0][0]=max(
    min_dp[0][0]+max(dp[v][0],dp[v][1]),
    min_dp[0][1]+dp[v][0]
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
   tmp[1][0]=max(
    min_dp[1][0]+max(dp[v][0],dp[v][1]),
    min_dp[1][1]+dp[v][0]
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
   memcpy(min_dp,tmp,sizeof tmp);
  dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
  dp[u][0]=min_dp[0][0];
int main(){
 int m,a,b;
 scanf("%d%d",&n,&m);
                                                                 }
 for(int i=0;i<m;i++){
  scanf("%d%d",&a,&b);</pre>
  init_g[a].push_back(b);
  init_g[b].push_back(a);
 par[1]=-1;
 tarjan(1);
 dfs(1,-1);
 printf("%d\n", max(dp[1][0], dp[1][1]));
 return 0;
8.7 Tree Knapsack
int dp[N][K]; PII obj[N];
vector<int> G[N];
void dfs(int u, int mx){
 for(int s: G[u]) {
  if(mx < obj[s].first) continue;</pre>
  for(int i=0;i<=mx-obj[s].FF;i++)</pre>
  dp[s][i] = dp[u][i];
dfs(s, mx - obj[s].first);
  for(int i=obj[s].FF;i<=mx;i++)</pre>
   dp[u][i] = max(dp[u][i],
    dp[s][i - obj[s].FF] + obj[s].SS);
}
8.8 N Queens Problem
vector< int > solve( int n ) {
 // no solution when n=2, 3
 vector< int > ret;
 if ( n % 6 == 2 ) {
 for ( int i = 2 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
  ret.push_back( 3 ); ret.push_back( 1 );
  for ( int i = 7 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 5 );
 } else if ( n % 6 == 3 ) {
for ( int i = 4 ; i <= n ; i += 2 )</pre>
   ret.push_back( i );
```

```
for ( int i = 5 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 1 ); ret.push_back( 3 );
 } else {
  for ( int i = 2 ; i <= n ; i += 2 )
   ret.push_back( i );
  for ( int i = 1 ; i <= n ; i += 2 )
   ret.push_back( i );
 return ret;
8.9 Aliens Optimization
long long Alien() {
 long long c = kInf;
 for (int d = 60; d >= 0; --d) {
  // cost can be negative, depending on the problem.
  if (c - (1LL << d) < 0) continue;</pre>
  long long ck = c - (1LL \ll d)
  pair<long long, int> r = check(ck);
  if (r.second == k) return r.first - ck * k;
  if (r.second < k) c = ck;
pair<long long, int> r = check(c);
 return r.first - c * k;
8.10 Hilbert Curve
long long hilbert(int n, int x, int y) {
 long long res = 0;
 for (int s = n / 2; s; s >>= 1) {
 int rx = (x & s) > 0, ry = (y & s) > 0;
res += s * 111 * s * ((3 * rx) ^ ry);
  if (ry == 0) {
   if (rx == 1) x = s - 1 - x, y = s - 1 - y;
   swap(x, y);
 return res;
8.11 Binary Search On Fraction
struct Q {
11 p, q;
 Q go(Q b, 11 d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(11 N) {
 Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
 assert(pred(hi));
 bool dir = 1, L = 1, H = 1;
 for (; L || H; dir = !dir) {
 11 len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)
   if (Q mid = hi.go(lo, len + step);
     mid.p > N \mid\mid mid.q > N \mid\mid dir ^ pred(mid))
    t++;
   else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
 return dir ? hi : lo;
```