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5	Mat	h	13	DKISEKI -Wall -Wextra -Wshadow -Wf
5	5.1		13	Wconversion -fsanitize=address,und
			13	success <cr></cr>
		5.2.1 First Kind		ap <f9> <esc>:w<cr>:!g++ "%" -o "%&lt;"</cr></esc></f9>
			13	echo success <cr></cr>
		3 3	13 13 m	ap <f10> <esc>:!./"%&lt;"<cr></cr></esc></f10>
			17	0 5 1 14
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			13   #	ifdef KISEKI
		3	14   #	define safe cerr< <pretty_function< td=""></pretty_function<>
			14 14	<<" line "< <line<<" safe\n"<="" td=""></line<<">
			14 #	<pre>define debug(a) qwerty(#a, a)</pre>
				define orange(a) dvorak(#a, a)
				sing std::cerr;
			45	emplate <typenamet></typenamet>
			15   v	<pre>roid qwerty(const char *s, Ta) {</pre>
			וסו	cerr << "\e[1;32m(" << s << ") = (";
			16 16	<pre>int cnt = sizeof(T);</pre>
			16	(, (cerr << a << (cnt ? ", " :
	5.20	Pollard Rho	16 }	•
		3 ,	16 t	emplate <typename iter=""></typename>
		~	17 v	oid dvorak(const char *s, Iter L, Ite
		•	17 18	cerr << "\e[1;32m[ " << s << " ] = [
	5.24	s simplex construction	10	for (int f = 0; L != R; ++L)
6	Geo	metry	18	cerr << (f++ ? ", " : "") << *L;
	6.1		18	cerr << " ]\e[0m\n";
			18 18 }	
			18   #	else
			19   #	define safe ((void)0)
	6.6	2D Closest Pair		<pre>define debug() ((void)0)</pre>
				define orange() ((void)0)
		5		endif
			19 20 <b>1</b>	7 Inches of Charles
			20 1	.3 Increase Stack
	6.12	Intersection of line and Circle	20 c	onst int size = 256 << 20;
		55		egister long rsp asm("rsp");
		3		har *p = (char*)malloc(size)+size, *b
			21 -	_asm("movq %0, %%rsp\n"::"r"(p));
			21 /	/ main
	6.17			(
			21 _	_asm("movq %0, %%rsp\n"::"r"(bak));

```
7 Stringology
           22
           22
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           23
           23
   23
   23
   23
   24
   24
   monotone (concave/convex) . . . . . . . . . . . . .
           24
   condition (concave/convex) . . . . . . . . . . . . . . .
           24
   24
   timization
           24
   25
           25
```

```
u mouse=a encoding=utf-8 ls=2
et sw=4 sts=4 tgc sc hls
ert
<CR>}<ESC>0
v<CR>:!g++ "%" -o "%<" -std=c++17 -
ll -Wextra -Wshadow -Wfatal-errors -
-fsanitize=address,undefined -g && echo
w<CR>:!g++ "%" -o "%<" -02 -std=c++17 &&</pre>
ss<CR>
!./"%<"<CR>
```

## lacro

```
rr<<__PRETTY_FUNCTION__\
_LINE__<<" safe\n"
...) qwerty(#a, a)
a...) dvorak(#a, a)
ame ...T>
st char *s, T ...a) {
;32m(" << s << ") = (";
eof...(T);
< a << (--cnt ? ", " : ")\e[0m\n")));</pre>
ame Iter>
st char *s, Iter L, Iter R) {
;32m[ " << s << " ] = [ ";
,
0;        L        != R; ++L)
+ ? ", " : "") << *L;
0m\n";
(0(biov
..) ((void)0)
...) ((void)0)
```

# Stack

```
= 256 << 20;
sp asm("rsp");
*)malloc(size)+size, *bak = (char*)rsp;
0, %%rsp\n"::"r"(p));
```

# 1.4 Pragma Optimization

```
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popent,abm,mmx,avx,tune=native")
1.5 IO Optimization
static inline int gc() {
constexpr int B = 1<<20;</pre>
 static char buf[B], *p, *q;
if(p == q \&\&
  (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
  return EOF:
return *p++;
template < typename T >
static inline bool gn( T &x ) {
int c = gc(); T sgn = 1; x = 0;
while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
if(c == '-') sgn = -1, c = gc();
if(c == EOF) return false;
while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
return x *= sgn, true;
```

#pragma GCC optimize("Ofast,no-stack-protector")

#pragma GCC optimize("no-math-errno,unroll-loops")

# 2 Data Structure

# 2.1 Dark Magic

#### 2.2 Link-Cut Tree

p->ch[dir]=c;

```
struct Node{
Node *par, *ch[2];
int xor_sum, v;
bool is_rev;
Node(int _v){
 v=xor_sum=_v;is_rev=false;
 par=ch[0]=ch[1]=nullptr;
inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
inline void down(){
 if(is_rev){
  if(ch[0]!=nullptr) ch[0]->set_rev();
   if(ch[1]!=nullptr) ch[1]->set_rev();
   is_rev=false;
 }
inline void up(){
 xor_sum=v;
  if(ch[0]!=nullptr){
  xor_sum^=ch[0]->xor_sum;
  ch[0]->par=this;
 if(ch[1]!=nullptr){
  xor_sum^=ch[1]->xor_sum;
  ch[1]->par=this;
 }
inline bool is_root(){
 return par==nullptr ||\
   (par->ch[0]!=this && par->ch[1]!=this);
bool is_rch(){return !is_root() && par->ch[1]==this;}
} *node[maxn], *stk[maxn];
int top;
void to_child(Node* p,Node* c,bool dir){
```

};

```
p->up();
inline void rotate(Node* node){
 Node* par=node->par;
 Node* par_par=par->par;
 bool dir=node->is_rch();
 bool par_dir=par->is_rch();
 to_child(par, node->ch[!dir], dir);
 to_child(node,par,!dir);
 if(par_par!=nullptr && par_par->ch[par_dir]==par)
  to_child(par_par,node,par_dir);
 else node->par=par_par;
inline void splay(Node* node){
 Node* tmp=node;
 stk[top++]=node;
 while(!tmp->is_root()){
  tmp=tmp->par;
  stk[top++]=tmp;
 while(top) stk[--top]->down();
 for(Node *fa=node->par;
  !node->is_root();
  rotate(node),fa=node->par)
  if(!fa->is_root())
   rotate(fa->is_rch()==node->is_rch()?fa:node);
inline void access(Node* node){
 Node* last=nullptr;
 while(node!=nullptr){
  splay(node);
  to_child(node, last, true);
  last=node;
  node=node->par;
inline void change_root(Node* node){
 access(node);splay(node);node->set_rev();
inline void link(Node* x, Node* y){
 change_root(x);splay(x);x->par=y;
inline void split(Node* x, Node* y){
 change_root(x);access(y);splay(x);
 to_child(x,nullptr,true);y->par=nullptr;
inline void change_val(Node* node,int v){
access(node);splay(node);node->v=v;node->up();
inline int query(Node* x,Node* y){
 change_root(x);access(y);splay(y);
 return y->xor_sum;
inline Node* find_root(Node* node){
 access(node);splay(node);
 Node* last=nullptr:
 while(node!=nullptr){
  node->down();last=node;node=node->ch[0];
 return last;
set<pii> dic;
inline void add_edge(int u,int v){
 if(u>v) swap(u,v)
 if(find_root(node[u])==find_root(node[v])) return;
 dic.insert(pii(u,v))
 link(node[u],node[v]);
inline void del_edge(int u,int v){
 if(u>v) swap(u,v);
 if(dic.find(pii(u,v))==dic.end()) return;
 dic.erase(pii(u,v))
 split(node[u],node[v]);
2.3 LiChao Segment Tree
struct L {
 int m, k, id;
 L(): id(-1) {}
L(int a, int b, int c) : m(a), k(b), id(c) {} int at(int x) { return m * x + k; }
```

```
class LiChao {
                                                             template <int BITS>
                                                             struct LinearBasis {
private:
                                                              array<uint64_t, BITS> basis;
int n; vector<L> nodes;
static int lc(int x) { return 2 * x + 1;
                                                              Basis() { basis.fill(0); }
static int rc(int x) { return 2 * x + 2; }
                                                              void add(uint64_t x) {
                                                               for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
void insert(int 1, int r, int id, L ln) {
                                                                if (basis[i] == 0) {
  int m = (1 + r) >> 1;
 if (nodes[id].id == -1) {
                                                                 basis[i] = x;
  nodes[id] = ln;
                                                                 return:
  return:
                                                                x ^= basis[i];
 bool atLeft = nodes[id].at(1) < ln.at(1);</pre>
                                                               }
 if (nodes[id].at(m) < ln.at(m)) {</pre>
  atLeft ^= 1
                                                              bool ok(uint64_t x) {
   swap(nodes[id], ln);
                                                               for (int i = 0; i < BITS; ++i)</pre>
                                                                if ((x >> i) & 1) x ^= basis[i];
  if (r - l == 1) return;
                                                               return x == 0;
  if (atLeft) insert(1, m, lc(id), ln);
                                                             };
 else insert(m, r, rc(id), ln);
                                                                   Binary Search On Segment Tree
int query(int 1, int r, int id, int x) {
 int ret = 0, m = (1 + r) >> 1;
if (nodes[id].id != -1)
                                                             // find_first = x -> minimal x s.t. check( [a, x) )
                                                             // find_last = x \rightarrow maximal x s.t. check([x, b))
  ret = nodes[id].at(x);
                                                             template <typename C>
  if (r - 1 == 1) return ret;
                                                             int find_first(int 1, const C &check) {
 if (x < m) return max(ret, query(1, m, lc(id), x));
return max(ret, query(m, r, rc(id), x));</pre>
                                                              if (1 >= n) return n;
                                                              1 += sz:
                                                              for (int i = height; i > 0; i--)
                                                               propagate(1 >> i);
public:
                                                              Monoid sum = identity;
LiChao(int n_{-}) : n(n_{-}), nodes(n * 4) {}
                                                              do {
void insert(L ln) { insert(0, n, 0, ln); }
                                                               while ((1 & 1) == 0) 1 >>= 1;
int query(int x) { return query(0, n, 0, x); }
                                                               if (check(f(sum, data[1]))) {
                                                                while (1 < sz) {
                                                                 propagate(1);
2.4 Treap
                                                                 1 <<= 1;
                                                                 auto nxt = f(sum, data[1]);
namespace Treap{
                                                                 if (not check(nxt)) {
#define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                                  sum = nxt;
struct node{
                                                                  1++:
 int size;
                                                                 }
 uint32_t pri;
                                                                }
 node *lc, *rc, *pa;
                                                                return 1 + 1 - sz;
 node():size(0),pri(rand()),lc(0),rc(0),pa(0)\{\}
 void pull() {
                                                               sum = f(sum, data[1++]);
  size = 1; pa = nullptr;
                                                              } while ((1 & -1) != 1);
  if ( lc ) { size += lc->size; lc->pa = this; }
                                                              return n;
   if ( rc ) { size += rc->size; rc->pa = this; }
 }
                                                             template <typename C>
};
                                                             int find_last(int r, const C &check) {
node* merge( node* L, node* R ) {
                                                              if (r <= 0) return -1;</pre>
 if ( not L or not R ) return L ? L : R;
                                                              r += sz;
  if ( L->pri > R->pri ) {
                                                              for (int i = height; i > 0; i--)
  L->rc = merge( L->rc, R ); L->pull();
                                                               propagate((r - 1) >> i);
  return L;
                                                              Monoid sum = identity;
  } else {
                                                              do {
  R->lc = merge( L, R->lc ); R->pull();
   return R;
                                                               while (r > 1 and (r & 1)) r >>= 1;
                                                               if (check(f(data[r], sum))) {
                                                                while (r < sz) {</pre>
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                                 propagate(r);
 if ( not rt ) L = R = nullptr;
                                                                 r = (r << 1) + 1;
 else if( sz( rt->lc ) + 1 <= k ) {
                                                                 auto nxt = f(data[r], sum);
                                                                 if (not check(nxt)) {
  split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                  sum = nxt;
  L->pull();
                                                                  r--;
  } else {
                                                                 }
  R = rt;
                                                                }
   split_by_size( rt->lc, k, L, R->lc );
                                                                return r - sz;
   R->pull();
  }
                                                               sum = f(data[r], sum);
                                                              } while ((r & -r) != r);
int getRank(node *o) {
                                                              return -1;
 int r = sz(o->lc)
  for (;o->pa != nullptr; o = o->pa)
   if (o->pa->rc != o) r += sz(o->pa->lc);
                                                             3
                                                                  Graph
  return r;
                                                             3.1 2-SAT (SCC)
#undef sz
                                                             class TwoSat{
```

private:

int n;

## 2.5 Linear Basis

```
vector<vector<int>> rG,G,sccs;
                                                                 void add_edge(int u, int v) {
                                                                 G[u].emplace_back(v, ecnt);
  vector<int> ord,idx;
  vector<bool> vis,result;
                                                                  G[v].emplace_back(u, ecnt++);
  void dfs(int u){
  vis[u]=true
                                                                 void solve() {
                                                                 bridge.assign(ecnt, false);
   for(int v:G[u])
    if(!vis[v]) dfs(v);
                                                                  for (int i = 0; i < n; ++i)
                                                                  if (not dfn[i]) dfs(i, i);
   ord.push_back(u);
  void rdfs(int u){
                                                                 bool is_bridge(int x) { return bridge[x]; }
  vis[u]=false;idx[u]=sccs.size()-1;
                                                              } bcc_bridge;
   sccs.back().push_back(u);
                                                              3.3 BCC Vertex
   for(int v:rG[u])
    if(vis[v])rdfs(v);
                                                              class BCC_AP {
                                                               private:
public:
                                                                int n, ecnt;
  void init(int n_){
                                                                 vector<vector<pair<int,int>>> G;
                                                                 vector<int> bcc, dfn, low, st;
  n=n_;G.clear();G.resize(n);
   rG.clear();rG.resize(n);
                                                                 vector<bool> ap, ins;
   sccs.clear();ord.clear();
                                                                 void dfs(int u, int f) {
                                                                  dfn[u] = low[u] = dfn[f] + 1;
   idx.resize(n);result.resize(n);
                                                                  int ch = 0
                                                                  for (auto [v, t]: G[u]) if (v != f) {
  if (not ins[t]) {
  void add_edge(int u,int v){
  G[u].push_back(v);rG[v].push_back(u);
                                                                    st.push_back(t);
  void orr(int x,int y){
                                                                    ins[t] = true;
  if ((x^y)==1)return
   add_edge(x^1,y); add_edge(y^1,x);
                                                                   if (dfn[v]) {
                                                                    low[u] = min(low[u], dfn[v]);
  bool solve(){
                                                                   } ++ch; dfs(v, u);
   vis.clear();vis.resize(n);
   for(int i=0;i<n;++i)</pre>
                                                                   low[u] = min(low[u], low[v]);
    if(not vis[i])dfs(i);
                                                                   if (low[v] >= dfn[u]) {
   reverse(ord.begin(),ord.end());
                                                                    ap[u] = true;
   for (int u:ord){
                                                                    while (true) {
    if(!vis[u])continue;
                                                                     int eid = st.back(); st.pop_back();
    sccs.push_back(vector<int>());
                                                                     bcc[eid] = ecnt;
    rdfs(u);
                                                                     if (eid == t) break;
   for(int i=0;i<n;i+=2)</pre>
                                                                    ecnt++:
    if(idx[i]==idx[i+1])
     return false;
   vector<bool> c(sccs.size());
                                                                  if (ch == 1 \text{ and } u == f) \text{ ap}[u] = false;
   for(size_t i=0;i<sccs.size();++i){</pre>
    for(size_t j=0;j<sccs[i].size();++j){
  result[sccs[i][j]]=c[i];</pre>
                                                               public:
                                                                 void init(int n_) {
     c[idx[sccs[i][j]^1]]=!c[i];
                                                                  G.clear(); G.resize(n = n_);
                                                                  ecnt = 0; ap.assign(n, false);
                                                                  low.assign(n, 0); dfn.assign(n, 0);
   return true;
                                                                 void add_edge(int u, int v) {
  bool get(int x){return result[x];}
                                                                 G[u].emplace_back(v, ecnt);
                                                                 G[v].emplace_back(u, ecnt++);
  inline int get_id(int x){return idx[x];}
  inline int count(){return sccs.size();}
} sat2;
                                                                 void solve() {
                                                                  ins.assign(ecnt, false);
3.2 BCC Edge
                                                                  bcc.resize(ecnt); ecnt = 0;
                                                                  for (int i = 0; i < n; ++i)
  if (not dfn[i]) dfs(i, i);</pre>
class BCC_Bridge {
private:
 int n, ecnt;
 vector<vector<pair<int,int>>> G;
                                                                 int get_id(int x) { return bcc[x]; }
  vector<int> dfn, low;
                                                                 int count() { return ecnt; }
  vector<bool> bridge;
                                                                 bool is_ap(int x) { return ap[x]; }
 void dfs(int u, int f) {
  dfn[u] = low[u] = dfn[f] + 1;
                                                              } bcc_ap;
                                                              3.4 Centroid Decomposition
   for (auto [v, t]: G[u]) {
    if (v == f) continue;
                                                              struct Centroid {
    if (dfn[v]) {
                                                               vector<vector<int64_t>> Dist;
                                                               vector<int> Parent, Depth;
     low[u] = min(low[u], dfn[v]);
     continue;
                                                               vector<int64_t> Sub, Sub2;
                                                               vector<int> Sz, Sz2
    dfs(v, u);
                                                               Centroid(vector<vector<pair<int, int>>> g) {
    low[u] = min(low[u], low[v]);
                                                                 int N = g.size();
    if (low[v] > dfn[u]) bridge[t] = true;
                                                                 vector<bool> Vis(N);
                                                                 vector<int> sz(N), mx(N);
                                                                 vector<int> Path;
public:
                                                                 Dist.resize(N)
  void init(int n_) {
                                                                 Parent.resize(N);
   G.clear(); G.resize(n = n_);
                                                                 Depth.resize(N)
                                                                auto DfsSz = [&](auto dfs, int x) -> void {
   Vis[x] = true; sz[x] = 1; mx[x] = 0;
   low.assign(n, ecnt = 0);
   dfn.assign(n, 0);
```

for (auto [u, w] : g[x]) {

```
if (Vis[u]) continue;
    dfs(dfs, u);
    sz[x] += sz[u]
    mx[x] = max(mx[x], sz[u]);
   Path.push_back(x);
  auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
   Dist[x].push_back(D); Vis[x] = true;
   for (auto [u, w] : g[x]) {
    if (Vis[u]) continue;
    dfs(dfs, u, D + w);
  };
  auto Dfs = [&]
   (auto dfs, int x, int D = 0, int p = -1)->void \{
   Path.clear(); DfsSz(DfsSz, x);
   int M = Path.size();
   int C = -1;
   for (int u : Path) {
    if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
    Vis[u] = false;
   DfsDist(DfsDist, C);
   for (int u : Path) Vis[u] = false;
   Parent[C] = p; Vis[C] = true;
   Depth[C] = D;
  for (auto [u, w] : g[C]) {
  if (Vis[u]) continue;
    dfs(dfs, u, D + 1, C);
   }
  Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
 Sz.resize(N); Sz2.resize(N);
void Mark(int v) {
  int x = v, z = -1;
  for (int i = Depth[v]; i >= 0; --i) {
  Sub[x] += Dist[v][i]; Sz[x]++;
   if (z != -1)
    Sub2[z] += Dist[v][i];
    Sz2[z]++;
   z = x; x = Parent[x];
  }
int64_t Query(int v) {
 int64_t res = 0;
 int x = v, z = -1;
 for (int i = Depth[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
  if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
  z = x; x = Parent[x];
  return res:
}
};
3.5 Directed Minimum Spanning Tree
```

```
struct DirectedMST { // find maximum
struct Edge {
 int u, v;
 int w:
 Edge(int u, int v, int w) : u(u), v(v), w(w) {}
}:
vector<Edge> Edges;
void clear() { Edges.clear(); }
void addEdge(int a, int b, int w) { Edges.emplace_back
    (a, b, w); }
int solve(int root, int n) {
 vector<Edge> E = Edges;
 int ans = 0;
 while (true) {
  // find best in edge
  vector<int> in(n, -inf), prv(n, -1);
  for (auto e : E)
   if (e.u != e.v && e.w > in[e.v]) {
     in[e.v] = e.w;
     prv[e.v] = e.u;
  in[root] = 0;
```

```
prv[root] = -1;
   for (int i = 0; i < n; i++)
    if (in[i] == -inf)
     return -inf;
   // find cycle
   int tot = 0;
   vector<int> id(n, -1), vis(n, -1);
for (int i = 0; i < n; i++) {</pre>
    ans += in[i];
    for (int x = i; x != -1 && id[x] == -1; x = prv[x])
     if (vis[x] == i) {
      for (int y = prv[x]; y != x; y = prv[y])
       id[y] = tot;
      id[x] = tot++;
      break;
     vis[x] = i;
    }
   if (!tot)
    return ans;
   for (int i = 0; i < n; i++)</pre>
    if (id[i] == -1)
     id[i] = tot++;
   // shrink
   for (auto &e : E) {
    if (id[e.u] != id[e.v])
     e.w -= in[e.v];
    e.u = id[e.u], e.v = id[e.v];
   n = tot:
   root = id[root];
  assert(false);
} DMST;
3.6 Dominator Tree
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
 // vertices are numbered from \theta to n-1
 fill(dfn, dfn + n, -1);fill(rev, rev + n, -1);
 fill(fa, fa + n, -1); fill(val, val + n, -1);
 fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
 fill(dom, dom + n, -1); tk = 0;
 for (int i = 0; i < n; ++i) {
  g[i].clear(); r[i].clear(); rdom[i].clear();
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
 rev[dfn[x] = tk] = x;
 fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
 for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
  r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
 if (fa[x] == x) return c ? -1 : x;
 int p = find(fa[x], 1);
 if (p == -1) return c ? fa[x] : val[x];
 if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
 fa[x] = p;
 return c ? p : val[x];
vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
// p[i] = -2 if i is unreachable from s
 dfs(s)
 for (int i = tk - 1; i >= 0; --i) {
  for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
  if (i) rdom[sdom[i]].push_back(i);
  for (int &u : rdom[i]) {
   int p = find(u);
   if (sdom[p] == i) dom[u] = i;
   else dom[u] = p;
```

```
if ( not chain[ u ] )
  if (i) merge(i, rp[i]);
                                                                  chain[ u ] = chain_ ++;
 }
vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)</pre>
                                                                void dfschain( int u, int f ) {
                                                                 fa[ u ][ 0 ] = f;
 if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
                                                                 for ( int i = 1 ; i < LOG_N ; ++ i )
                                                                  fa[u][i] = fa[fa[u][i-1]][i-1];
 for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
 return p;
                                                                 tl[ u ] = time_++
                                                                 if ( not chain_st[ chain[ u ] ] )
                                                                  chain_st[ chain[ u ] ] = u;
3.7 Edge Coloring
                                                                 for ( int v : G[ u ] )
                                                                  if ( v != f and chain[ v ] == chain[ u ] )
// \max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
                                                                   dfschain( v, u );
                                                                 for ( int v : G[ u ] )
void clear(int N) {
for (int i = 0; i <= N; i++)
for (int j = 0; j <= N; j++)
                                                                  if ( v != f and chain[ v ] != chain[ u ] )
                                                                   dfschain( v, u );
    C[i][j] = G[i][j] = 0;
                                                                 tr[ u ] = time_;
                                                                bool anc( int u, int v ) {
  return tl[ u ] <= tl[ v ] and tr[ v ] <= tr[ u ];
void solve(vector<pair<int, int>> &E, int N) {
 int X[kN] = {}, a;
 auto update = [&](int u)
                                                               public:
  for (X[u] = 1; C[u][X[u]]; X[u]++);
                                                                int lca( int u, int v ) {
 auto color = [&](int u, int v, int c) {
                                                                 if ( anc( u, v ) ) return u;
                                                                 for ( int i = LOG_N - 1; i >= 0; -- i )
if ( not anc( fa[ u ][ i ], v ) )
u = fa[ u ][ i ];
  int p = G[u][v];
  G[u][v] = G[v][u] = c;
  C[u][c] = v, C[v][c] = u;
  C[u][p] = C[v][p] = 0;
                                                                 return fa[ u ][ 0 ];
  if (p) X[u] = X[v] = p;
  else update(u), update(v);
                                                                void init( int n ) {
                                                                 fa.assign( ++n, vector< int >( LOG_N ) );
  return p;
                                                                 for ( LOG_N = 0 ; ( 1 << LOG_N ) < n ; ++ LOG_N );
 };
 auto flip = [&](int u, int c1, int c2) {
                                                                 G.clear(); G.resize( n );
 int p = C[u][c1];
                                                                 tl.assign( n, 0 ); tr.assign( n, 0 );
 swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
                                                                 chain.assig( n, 0 ); chain_st.assign( n, 0 );
  if (!C[u][c1]) X[u] = c1;
                                                                void add_edge( int u , int v ) {
  if (!C[u][c2]) X[u] = c2;
                                                                 // 1-base
  return p;
                                                                 G[ u ].push_back( v );
                                                                 G[v].push_back(u);
 for (int i = 1; i <= N; i++) X[i] = 1;
 for (int t = 0; t < E.size(); t++) {
  auto [u, v] = E[t];
                                                                void decompose(){
                                                                 chain_ = 1;
                                                                 predfs( 1, 1 );
  int v0 = v, c = X[u], c0 = c, d;
  vector<pair<int, int>> L; int vst[kN] = {};
                                                                 time_{-} = 0;
                                                                 dfschain(1,1);
  while (!G[u][v0]) {
   L.emplace_back(v, d = X[v]);
if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
                                                                PII get_subtree(int u) { return {tl[ u ],tr[ u ] }; }
                                                                vector< PII > get_path( int u , int v ){
     c = color(u, L[a].first, c);
   else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
                                                                 vector< PII > res;
                                                                 int g = lca( u, v );
while ( chain[ u ] != chain[ g ] ) {
     color(u, L[a].first, L[a].second);
   else if (vst[d]) break
                                                                  int s = chain_st[ chain[ u ] ];
   else vst[d] = 1, v = C[u][d];
                                                                  res.emplace_back( tl[ s ], tl[ u ] + 1 );
  if (!G[u][v0]) {
  for (; v; v = flip(v, c, d), swap(c, d));
                                                                  u = fa[ s ][ 0 ];
                                                                 res.emplace_back( tl[ g ], tl[ u ] + 1 );
while ( chain[ v ] != chain[ g ] ) {
   if (C[u][c0]) { a = int(L.size()) - 1;
    while (--a >= 0 && L[a].second != c);
                                                                  int s = chain_st[ chain[ v ] ];
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
                                                                  res.emplace_back( tl[s], tl[v] + 1);
   } else t--;
                                                                  v = fa[s][0];
                                                                 res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
                                                                 return res;
3.8 Lowbit Decomposition
                                                                 /* res : list of intervals from u to v
class LowbitDecomp{
                                                                  * ( note only nodes work, not edge )
                                                                  * usage :
private:
                                                                  * vector< PII >& path = tree.get_path( u , v )
 int time_, chain_, LOG_N;
                                                                  * for( auto [ 1, r ] : path ) {
 vector< vector< int > > G, fa;
 vector< int > tl, tr, chain, chain_st;
                                                                  * 0-base [ 1, r )
                                                                  * }
 // chain_ : number of chain
 // tl, tr[ u ] : subtree interval in the seq. of u
                                                                  */
                                                                }
 // chain_st[ u ] : head of the chain contains u
 // chian[ u ] : chain id of the chain u is on
                                                               } tree;
 void predfs( int u, int f ) {
                                                                    Manhattan Minimum Spanning Tree
  chain[ u ] = 0;
  for ( int v : G[ u ] ) {
  if ( v == f ) continue;
                                                               typedef Point<int> P;
                                                               vector<array<int, 3>> manhattanMST(vector<P> ps) {
   predfs( v, u );
                                                                vi id(sz(ps));
   if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
                                                                iota(all(id), 0);
    chain[ u ] = chain[ v ];
                                                                vector<array<int, 3>> edges;
                                                                rep(k, 0, 4) {
```

bits pob, nob = 0; pob.set();

```
sort(all(id), [&](int i, int j) {
                                                                     for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
   return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
                                                                     for ( size_t i = 0 ; i < n ; ++ i ) {
                                                                      size_t v = deo[ i ];
  }):
                                                                      bits tmp; tmp[ v ] = 1;
  map<int, int> sweep;
                                                                      BK( tmp, pob & G[ v ], nob & G[ v ] );
  for (int i : id) {
                                                                      pob[v] = 0, nob[v] = 1;
   for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++)) {
    int i = it->second
                                                                     return static_cast< int >( ans.count() );
    Pd = ps[i] - ps[j];
    if (d.y > d.x) break;
                                                                  };
    edges.push_back({d.y + d.x, i, j});
                                                                  3.11 MaxCliqueDyn
   sweep[-ps[i].y] = i;
                                                                  constexpr int kN = 150;
                                                                   struct MaxClique { // Maximum Clique
  for (P &p : ps)
                                                                    bitset<kN> a[kN], cs[kN];
int ans, sol[kN], q, cur[kN], d[kN], n;
   if (k \& 1) p.x = -p.x;
   else swap(p.x, p.y);
                                                                    void init(int _n) {
                                                                    n = _n, ans = q = 0;
for (int i = 0; i < n; i++) a[i].reset();</pre>
return edges; // [{w, i, j}, ...]
3.10 MaxClique
                                                                    void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
void csort(vector<int> &r, vector<int> &c) {
// contain a self loop u to u, than u won't in clique
                                                                     int mx = 1, km = max(ans - q + 1, 1), t = 0,
template < size_t MAXN >
                                                                       m = int(r.size())
class MaxClique{
                                                                     cs[1].reset(); cs[2].reset();
private:
                                                                     for (int i = 0; i < m; i++) {
using bits = bitset< MAXN >;
                                                                      int p = r[i], k = 1;
while ((cs[k] & a[p]).count()) k++;
bits popped, G[ MAXN ], ans;
size_t deg[ MAXN ], deo[ MAXN ], n;
                                                                      if (k > mx) cs[++mx + 1].reset();
void sort_by_degree() {
                                                                      cs[k][p] = 1;
  popped.reset();
                                                                      if (k < km) r[t++] = p;
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                     }
    deg[ i ] = G[ i ].count();
  for ( size_t i = 0 ; i < n ; ++ i ) {
                                                                     c.resize(m);
                                                                     if(t) c[t-1] = 0;
    size_t mi = MAXN, id = 0;
    for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )</pre>
                                                                     for (int k = km; k <= mx; k++) {</pre>
                                                                      for (int p = int(cs[k]._Find_first());
                                                                         p < kN; p = int(cs[k]._Find_next(p))) {
        mi = deg[id = j];
    popped[ deo[ i ] = id ] = 1;
for( size_t u = G[ i ]._Find_first() ;
                                                                       r[t] = p; c[t++] = k;
                                                                     }
     u < n ; u = G[ i ]._Find_next( u ) )
      -- deg[ u ];
                                                                    void dfs(vector<int> &r, vector<int> &c, int 1,
  }
                                                                     bitset<kN> mask) {
void BK( bits R, bits P, bits X ) {
  if (R.count()+P.count() <= ans.count()) return;</pre>
                                                                     while (!r.empty()) {
                                                                      int p = r.back(); r.pop_back();
                                                                      mask[p] = 0;
  if ( not P.count() and not X.count() ) {
                                                                      if (q + c.back() <= ans) return;</pre>
  if ( R.count() > ans.count() ) ans = R;
                                                                      cur[q++] = p;
   return;
                                                                      vector<int> nr, nc;
  }
                                                                      bitset<kN> nmask = mask & a[p];
  /* greedily chosse max degree as pivot
                                                                      for (int i : r)
  bits cur = P | X; size_t pivot = 0, sz = 0;
                                                                       if (a[p][i]) nr.push_back(i);
  for ( size_t u = cur._Find_first() ;
                                                                      if (!nr.empty()) {
   u < n ; u = cur._Find_next( u )</pre>
                                                                       if (1 < 4) {
  if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
cur = P & ( ~G[ pivot ] );
                                                                        for (int i : nr)
                                                                         d[i] = int((a[i] & nmask).count());
  */ // or simply choose first
                                                                        sort(nr.begin(), nr.end(),
  bits cur = P & (~G[ ( P | X )._Find_first() ]);
                                                                         [&](int x, int y)
  for ( size_t u = cur._Find_first()
                                                                           return d[x] > d[y];
   u < n ; u = cur._Find_next( u ) ) {
   if ( R[ u ] ) continue;
R[ u ] = 1;
                                                                         });
                                                                      csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
   BK( R, P & G[ u ], X & G[ u ]);
   R[u] = P[u] = 0, X[u] = 1;
                                                                       ans = q; copy(cur, cur + q, sol);
                                                                      c.pop_back(); q--;
public:
                                                                     }
void init( size_t n_ ) {
                                                                    int solve(bitset<kN> mask) { // vertex mask
  for ( size_t i = 0 ; i < n ; ++ i )
  G[ i ].reset();
                                                                     vector<int> r, c
                                                                     for (int i = 0; i < n; i++)</pre>
  ans.reset();
                                                                     if (mask[i]) r.push_back(i);
for (int i = 0; i < n; i++)</pre>
void add_edges( int u, bits S ) { G[ u ] = S; }
void add_edge( int u, int v ) {
  G[ u ][ v ] = G[ v ][ u ] = 1;
                                                                      d[i] = int((a[i] & mask).count());
                                                                     sort(r.begin(), r.end(),
  [&](int i, int j) { return d[i] > d[j]; });
                                                                     csort(r, c);
int solve() {
                                                                     dfs(r, c, 1, mask);
return ans; // sol[0 \sim ans-1]
  sort_by_degree(); // or simply iota( deo... )
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
   deg[ i ] = G[ i ].count()
```

} graph;

# 3.12 Minimum Mean Cycle

```
/* minimum mean cycle O(VE) */
struct MMC{
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
#define V 1021
#define inf 1e9
 struct Edge { int v,u; double c; };
 int n, m, prv[V][V], prve[V][V], vst[V];
 Edge e[E];
 vector<int> edgeID, cycle, rho;
 double d[V][V];
 void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
 void add_edge( int vi , int ui , double ci )
{ e[ m ++ ] = { vi , ui , ci }; }
 void bellman_ford() {
  for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);</pre>
   for(int j=0; j<m; j++) {
  int v = e[j].v, u = e[j].u;
  if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
      d[i+1][u] = d[i][v]+e[j].c;
      prv[i+1][u] = v;
      prve[i+1][u] = j;
   }
  }
 double solve(){
  // returns inf if no cycle, mmc otherwise
  double mmc=inf;
  int st = -1
  bellman_ford();
  for(int i=0; i<n; i++) {</pre>
   double avg=-inf;
   for(int k=0; k<n; k++) {</pre>
    if(d[n][i]<inf-eps]</pre>
      avg=max(avg,(d[n][i]-d[k][i])/(n-k));
    else avg=max(avg,inf);
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
  for (int i=n; !vst[st]; st=prv[i--][st]) {
   vst[st]++
   edgeID.PB(prve[i][st]);
   rho.PB(st);
  while (vst[st] != 2) {
   int v = rho.back(); rho.pop_back();
   cycle.PB(v);
   vst[v]++;
  reverse(ALL(edgeID));
  edgeID.resize(SZ(cycle));
  return mmc;
} mmc;
3.13 Minimum Steiner Tree
// Minimum Steiner Tree
```

```
// 0(V 3^T + V^2 2^T)
struct SteinerTree {
#define V 33
#define T 8
#define INF 1023456789
int n, dst[V][V], dp[1 << T][V], tdst[V];</pre>
void init(int _n) {
  for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++)
dst[i][j] = INF * (i != j);</pre>
void add_edge(int ui, int vi, int wi) {
 dst[ui][vi] = min(dst[ui][vi], wi);
 dst[vi][ui] = min(dst[vi][ui], wi);
void shortest_path() {
```

```
for (int k = 0; k < n; k++)
   for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < n; j++)
      dst[i][j] = min(dst[i][j], dst[i][k] + dst[k][j]);
 int solve(const vector<int> &ter) {
  int t = (int)ter.size();
  for (int i = 1; i < (1 << t); i++)
   fill_n(dp[i], n, INF);
  fill_n(dp[0], n, 0);
for (int msk = 1; msk < (1 << t); msk++) {
   if (msk == (msk & (-msk))) {
    int who = __lg(msk);
for (int i = 0; i < n; i++)</pre>
     dp[msk][i] = dst[ter[who]][i];
    continue;
   for (int i = 0; i < n; i++)
    for (int submsk = (msk - 1) & msk; submsk; submsk =
      (submsk - 1) & msk)
     dp[msk][i] = min(dp[msk][i], dp[submsk][i] + dp[
    msk ^ submsk][i]);
   for (int i = 0; i < n; i++) {
    tdst[i] = INF
    for (int j = 0; j < n; j++)
     tdst[i] = min(tdst[i], dp[msk][j] + dst[j][i]);
   copy_n(tdst, n, dp[msk]);
  }
  int ans = INF;
  for (int i = 0; i < n; i++)</pre>
   ans = min(ans, dp[(1 << t) - 1][i]);
  return ans;
} solver;
3.14 Mo's Algorithm on Tree
int q; vector< int > G[N];
struct Que{
 int u, v,
} que[ N ];
int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
void_dfs( int u, int f ) {
 dfn[ u ] = dfn_++; int saved_rbp = stk_;
 for ( int v : G[ u ] ) {
  if ( v == f ) continue;
  dfs( v, u );
  if ( stk_ - saved_rbp < SQRT_N ) continue;
for ( ++ block_ ; stk_ != saved_rbp ; )
  block_id[ stk[ -- stk_ ] ] = block_;</pre>
 stk[ stk_ ++ ] = u;
bool inPath[ N ];
void Diff( int u ) {
  if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
 else { /*add this edge*/ }
void traverse( int& origin_u, int u ) {
for ( int g = lca( origin_u, u ) ;
  origin_u != g ; origin_u = parent_of[ origin_u ] )
   Diff( origin_u );
 for (int v = u; v != origin_u; v = parent_of[v])
  Diff( v );
 origin_u = u;
}
void solve() {
 dfs( 1, 1 );
while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
 sort( que, que + q, [](const Que& x, const Que& y) {
  return tie( block_id[ x.u ], dfn[ x.v ] )
       < tie( block_id[ y.u ], dfn[ y.v ] );
 int U = 1, V = 1;
for ( int i = 0 ; i < q ; ++ i ) {
 pass( U, que[ i ].u );
  pass( V, que[ i ].v );
  // we could get our answer of que[ i ].id
Method 2:
```

for (int v : G[ u ]) if (v != f) {

r=(r+(hh\*hh)%1010101333)%1011820613;

uint64\_t hh = hsah(v, u);

```
dfs u:
  push u
  iterate subtree
  push u
Let P = LCA(u, v), and St(u)<=St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
*/
3.15  Tree Hashing
uint64_t hsah(int u, int f) {
  uint64_t r = 127;</pre>
```

# 3.16 Virtural Tree

return r;

```
inline bool cmp(const int &i, const int &j) {
return dfn[i] < dfn[j];</pre>
void build(int vectrices[], int k) {
static int stk[MAX_N];
sort(vectrices, vectrices + k, cmp);
stk[sz++] = 0;
for (int i = 0; i < k; ++i)
 int u = vectrices[i], lca = LCA(u, stk[sz - 1]);
 if (lca == stk[sz - 1]) stk[sz++] = u;
 else {
  while (sz \ge 2 \&\& dep[stk[sz - 2]] \ge dep[lca]) {
   addEdge(stk[sz - 2], stk[sz - 1]);
  if (stk[sz - 1] != lca) {
   addEdge(lca, stk[--sz]);
    stk[sz++] = lca, vectrices[cnt++] = lca;
  stk[sz++] = u;
 }
for (int i = 0; i < sz - 1; ++i)
 addEdge(stk[i], stk[i + 1]);
```

# 4 Matching & Flow

# 4.1 Bipartite Matching

```
struct BipartiteMatching {
vector<int> X[N];
int fX[N], fY[N], n;
bitset<N> vis;
bool dfs(int x)
 for (auto i:X[x]) {
  if (vis[i]) continue;
   vis[i] = true;
  if (fY[i]==-1 || dfs(fY[i])){
    fY[fX[x] = i] = x;
    return true:
   }
  return false;
void init(int n_, int m) {
 vis.reset();
 fill(X, X + (n = n_{-}), vector<int>());
 memset(fX, -1, sizeof(int) * n);
 memset(fY, -1, sizeof(int) * m);
void add_edge(int x, int y){
 X[x].push_back(y); }
 int solve() { // return how many pair matched
 int cnt = 0:
 for(int i=0;i<n;i++) {</pre>
  vis.reset();
   cnt += dfs(i);
  }
  return cnt;
};
```

# 4.2 Dijkstra Cost Flow

```
// kN = #(vertices)
// MCMF.{Init, AddEdge, MincostMaxflow}
// MincostMaxflow(source, sink, flow_limit, &cost)
// => flow
using Pii = pair<int, int>
constexpr int kInf = 0x3f3f3f3f, kN = 500;
struct Edge {
 int to, rev, cost, flow;
struct MCMF { // 0-based
 int n{}, m{}, s{}, t{};
 vector<Edge> graph[kN];
 // Larger range for relabeling
 int64_t dis[kN] = {}, h[kN] = {};
 int p[kN] = {};
 void Init(int nn) {
  n = nn;
  for (int i = 0; i < n; i++) graph[i].clear();</pre>
 void AddEdge(int u, int v, int f, int c) {
  graph[u].push_back({v,
   static_cast<int>(graph[v].size()), c, f});
  graph[v].push_back(
   {u, static_cast<int>(graph[u].size()) - 1,
     -c, 0});
 bool Dijkstra(int &max_flow, int64_t &cost) {
  priority_queue<Pii, vector<Pii>, greater<>> pq;
  fill_n(dis, n, kInf);
  dis[s] = 0
  pq.emplace(0, s)
  while (!pq.empty()) {
   auto u = pq.top();
   pq.pop();
   int v = u.second;
   if (dis[v] < u.first) continue;</pre>
   for (auto &e : graph[v]) {
    auto new_dis =
     dis[v] + e.cost + h[v] - h[e.to];
    if (e.flow > 0 && dis[e.to] > new_dis) {
     dis[e.to] = new_dis;
     p[e.to] = e.rev;
     pq.emplace(dis[e.to], e.to);
    }
  if (dis[t] == kInf) return false;
  for (int i = 0; i < n; i++) h[i] += dis[i];</pre>
  int d = max_flow;
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   d = min(d, graph[e.to][e.rev].flow);
  max_flow -= d;
  cost += int64_t(d) * h[t];
  for (int u = t; u != s
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   e.flow += d;
   graph[e.to][e.rev].flow -= d;
  }
  return true;
 int MincostMaxflow(
  int ss, int tt, int max_flow, int64_t &cost) {
  this->s = ss, this->t = tt;
  cost = 0;
  fill_n(h, n, 0);
  auto orig_max_flow = max_flow;
  while (Dijkstra(max_flow, cost) && max_flow) {}
  return orig_max_flow - max_flow;
};
4.3 Dinic
```

```
template <typename Cap = int64_t>
class Dinic{
private:
    struct E{
```

```
int to, rev;
    Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
     queue<int> bfs;
    bfs.push(st); lv[st] = 0;
while (not bfs.empty()){
       int u = bfs.front(); bfs.pop();
       for (auto e: G[u]) {
         if (e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
         bfs.push(e.to); lv[e.to] = lv[u] + 1;
     return lv[ed] != -1;
  Cap DFS(int u, Cap f){
     if (u == ed) return f;
     Cap ret = 0;
     for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
       auto &e = G[u][i];
       if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
       Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
       if (f == 0) return ret;
     if (ret == 0) lv[u] = -1;
     return ret:
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
     G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
     st = st_, ed = ed_; Cap ret = 0;
    while (BFS()) {
       idx.assign(n, 0);
       Cap f = DFS(st, numeric_limits<Cap>::max());
       ret += f;
       if (f == 0) break;
     return ret:
};
```

## Flow Models

- · Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.

  - 2. For each edge (x,y,l,u), connect  $x\to y$  with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v \to T$  with capacity -in(v).
    - To maximize, connect  $t\to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T.If  $f 
      eq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the
    - maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect t o s with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v\in V,\, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e + f_e$ , where  $f_e$  corresponds to the flow of edge  $\boldsymbol{e}$  on the graph.
- Construct minimum vertex cover from maximum matchina M on bipartite graph(X,Y)
  - 1. Redirect every edge:  $y \to x$  if  $(x,y) \in M$ ,  $x \to y$  otherwise. 2. DFS from unmatched vertices in X.

  - 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source  ${\cal S}$  and sink  ${\cal T}$
  - 2. For each edge (x,y,c), connect  $x \to y$  with (cost,cap) = (c,1) if
  - c>0, otherwise connect  $y\to x$  with (cost, cap)=(c,1) if c>0, otherwise connect  $y\to x$  with (cost, cap)=(-c,1) 3. For each edge with c<0, sum these cost as K, then increase d(y)by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v)>0, connect  $S\to v$  with (cost, cap)=
  - 5. For each vertex v with d(v) < 0, connect  $v \to T$  with (cost, cap) =(0, -d(v))

- 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let K be the sum of all weights
  - 3. Connect source  $s \to v$ ,  $v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity
  - 5. For  $v \in {\it G}$ , connect it with sink  $v \to t$  with capacity K+2T- $(\sum_{e \in E(v)} w(e)) - 2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight w(u,v)
  - 2. Connect v o v' with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 3. Find the minimum weight perfect matching on  $G^{\prime}$ .
- · Project selection problem
  - 1. If  $p_v>0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge
  - (v,t) with capacity  $-p_v$ . 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
  - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

namespace matching {

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity  $c_x$  and create edge (s,y) with ca-
- 2. Create edge (x,y) with capacity  $c_{xy}$ . 3. Create edge (x,y) and edge (x',y') with capacity  $c_{xyx'y'}$ .

# 4.5 General Graph Matching

```
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
g[u].push_back(v);
 g[v].push_back(u);
int Find(int u) {
 return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
 static int tk = 0; tk++;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
v[x] = tk;
   x = Find(pre[match[x]]);
  }
void Blossom(int x, int y, int 1) {
 while (Find(x) != 1) {
  pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
  if (fa[x] == x) fa[x] = 1;
  if (fa[y] == y) fa[y] = 1;
  x = pre[y];
bool Bfs(int r, int n) {
 for (int i = 0; i \le n; ++i) fa[i] = i, s[i] = -1;
 while (!q.empty()) q.pop();
 q.push(r);
 s[r] = 0;
 while (!q.empty()) {
  int x = q.front(); q.pop();
  for (int u : g[x]) {
  if (s[u] == -1) {
    pre[u] = x, s[u] = 1;
    if (match[u] == n) {
     for (int a = u, b = x, last; b != n; a = last, b =
     pre[a])
```

```
last = match[b], match[b] = a, match[a] = b;
    return true;
}
    q.push(match[u]);
    s[match[u]] = 0;
} else if (!s[u] && Find(u) != Find(x)) {
    int l = LCA(u, x, n);
    Blossom(x, u, l);
    Blossom(u, x, l);
}

return false;
}
int Solve(int n) {
    int res = 0;
    for (int x = 0; x < n; ++x) {
        if (match[x] == n) res += Bfs(x, n);
    }
    return res;
}</pre>
```

#### 4.6 Global Min-Cut

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
pair<int, int> phase(int n) {
memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
int s = -1, t = -1;
while (true) {
 int c = -1;
 for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
   if (c == -1 \mid | g[i] > g[c]) c = i;
  if (c == -1) break;
 v[s = t, t = c] = true;
 for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
   g[i] += w[c][i];
 }
 return make_pair(s, t);
int mincut(int n) {
int cut = 1e9;
memset(del, false, sizeof(del));
for (int i = 0; i < n - 1; ++i) {
  int s, t; tie(s, t) = phase(n);
 del[t] = true; cut = min(cut, g[t]);
 for (int j = 0; j < n; ++j) {
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
 }
return cut;
```

### 4.7 GomoryHu Tree

```
int g[maxn];
vector<edge> GomoryHu(int n){
  vector<edge> rt;
  for(int i=1;i<=n;++i)g[i]=1;
  for(int i=2;i<=n;++i){
    int t=g[i];
    flow.reset(); // clear flows on all edge
    rt.push_back({i,t,flow(i,t)});
    flow.walk(i); // bfs points that connected to i (use
        edges not fully flow)
    for(int j=i+1;j<=n;++j){
        if(g[j]==t && flow.connect(j))g[j]=i; // check if i
            can reach j
    }
}
return rt;
}</pre>
```

#### 4.8 Kuhn Munkres

```
class KM {
private:
 static constexpr lld INF = 1LL << 60;</pre>
 vector<lld> hl,hr,slk;
 vector<int> fl,fr,pre,qu;
 vector<vector<lld>> w;
 vector<bool> v1, vr;
 int n, ql, qr;
 bool check(int x) {
  if (v1[x] = true, f1[x] != -1)
   return vr[qu[qr++] = f1[x]] = true;
  while (x != -1) swap(x, fr[fl[x] = pre[x]]);
  return false;
 void bfs(int s) {
  fill(slk.begin(), slk.end(), INF);
  fill(v1.begin(), v1.end(), false);
fill(vr.begin(), vr.end(), false);
  ql = qr = 0;
  vr[qu[qr++] = s] = true;
  while (true) {
   11d d;
    while (ql < qr) {</pre>
     for (int x = 0, y = qu[ql++]; x < n; ++x) {
   if(!v1[x]&&s1k[x]>=(d=h1[x]+hr[y]-w[x][y])){
       if (pre[x] = y, d) slk[x] = d;
       else if (!check(x)) return;
     }
    }
    d = INF;
   for (int x = 0; x < n; ++x)
if (!v1[x] && d > s1k[x]) d = s1k[x];
    for (int x = 0; x < n; ++x) {
    if (v1[x]) h1[x] += d;
     else slk[x] -= d;
     if (vr[x]) hr[x] -= d;
    for (int x = 0; x < n; ++x)
     if (!v1[x] && !slk[x] && !check(x)) return;
public:
 void init( int n_ ) {
  qu.resize(n = n_);
  fl.assign(n, -1); fr.assign(n, -1);
hr.assign(n, 0); hl.resize(n);
  w.assign(n, vector<lld>(n));
  slk.resize(n); pre.resize(n);
  vl.resize(n); vr.resize(n);
 void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 11d solve() {
  for (int i = 0; i < n; ++i)
   hl[i] = *max_element(w[i].begin(), w[i].end());
  for (int i = 0; i < n; ++i) bfs(i);</pre>
  11d res = 0;
  for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
  return res;
 }
} km;
```

# 4.9 Minimum Cost Circulation

eg.cap -= mw;

```
while(!mark[upd])mark[upd]=1,upd=pv[upd];
                                                               G[eg.to][eg.back].cap+=mw;
      return upd:
     }
                                                              return {mw, dis[edd]};
    idx++:
                                                            public:
                                                             void init(int n){
 }
                                                              G.clear();G.resize(n);
                                                              fa.resize(n);wh.resize(n);
return -1;
                                                              inq.resize(n); dis.resize(n);
int Solve(int n) {
                                                             void add_edge(int st, int ed, Cap c, Wei w){
int rt = -1, ans = 0;
                                                              G[st].emplace_back(ed,SZ(G[ed]),c,w);
while ((rt = NegativeCycle(n)) >= 0) {
                                                              G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
 memset(mark, false, sizeof(mark));
 vector<pair<int, int>> cyc;
                                                             PCW solve(int a, int b){
                                                              ori = a, edd = b;
Cap cc=0; Wei ww=0;
 while (!mark[rt]) {
  cyc.emplace_back(pv[rt], ed[rt]);
  mark[rt] = true;
                                                              while(true){
                                                               PCW ret=SPFA();
  rt = pv[rt];
                                                               if(ret.first==-1) break;
 reverse(cyc.begin(), cyc.end());
                                                               cc+=ret.first:
                                                               ww+=ret.first * ret.second;
 int cap = kInf;
 for (auto &i : cyc)
  auto &e = g[i.first][i.second];
                                                              return {cc,ww};
  cap = min(cap, e.cap);
                                                            } mcmf;
 for (auto &i : cyc) {
                                                            4.11 Minimum Weight Matching (Clique version)
  auto &e = g[i.first][i.second];
  e.cap -= cap;
                                                            struct Graph {
                                                             // 0-base (Perfect Match)
  g[e.to][e.rev].cap += cap;
  ans += e.cost * cap;
                                                             int n, edge[MXN][MXN];
 }
                                                             int match[MXN], dis[MXN], onstk[MXN];
                                                             vector<int> stk;
return ans;
                                                             void init(int _n) {
                                                              n = _n;
                                                              for (int i=0; i<n; i++) for (int j=0; j<n; j++)
4.10
     Minimum Cost Maximum Flow
                                                               edge[i][j] = 0;
class MiniCostMaxiFlow{
using Cap = int; using Wei = int64_t;
                                                             void set_edge(int u, int v, int w) {
                                                              edge[u][v] = edge[v][u] = w; }
using PCW = pair<Cap,Wei>;
static constexpr Cap INF_CAP = 1 << 30;</pre>
                                                             bool SPFA(int u){
static constexpr Wei INF_WEI = 1LL<<60;</pre>
                                                              if (onstk[u]) return true;
private:
                                                              stk.PB(u); onstk[u] = 1;
struct Edge{
                                                              for (int v=0; v<n; v++){
 int to, back;
Cap cap; Wei wei;
                                                               if (u != v && match[u] != v && !onstk[v]){
                                                                int m = match[v]
 Edge() {}
                                                                if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
 Edge(int a,int b, Cap c, Wei d):
                                                                 dis[m] = dis[u] - edge[v][m] + edge[u][v];
                                                                 onstk[v] = 1;
   to(a),back(b),cap(c),wei(d) {}
};
                                                                 stk.PB(v);
int ori, edd;
                                                                 if (SPFA(m)) return true;
vector<vector<Edge>> G;
                                                                 stk.pop_back();
vector<int> fa, wh;
                                                                 onstk[v] = 0;
vector<bool> inq;
vector<Wei> dis;
PCW SPFA(){
 fill(inq.begin(),inq.end(),false);
                                                              onstk[u] = 0; stk.pop_back();
 fill(dis.begin(), dis.end(), INF_WEI);
                                                              return false:
  queue<int> qq; qq.push(ori);
 dis[ori] = 0;
                                                             int solve() { // find a match
 while(not qq.empty()){
                                                              for (int i=0; i<n; i+=2){</pre>
                                                               match[i] = i+1;
  int u=qq.front();qq.pop();
  ing[u] = false;
                                                               match[i+1] = i;
   for(int i=0;i<SZ(G[u]);++i){</pre>
   Edge e=G[u][i];
                                                              while (true){
    int v=e.to; Wei d=e.wei;
                                                               int found = 0;
                                                               for (int i=0; i<n; i++)</pre>
    if(e.cap <= 0 | |dis[v] <= dis[u] + d)
                                                                dis[i] = onstk[i] = 0;
     continue
    dis[v] = dis[u] + d;
                                                               for (int i=0; i<n; i++){</pre>
   fa[v] = u, wh[v] = i;
                                                                stk.clear()
                                                                if (!onstk[i] && SPFA(i)){
   if (ing[v]) continue;
                                                                 found = 1
    qq.push(v);
    inq[v] = true;
                                                                 while (SZ(stk)>=2){
  }
                                                                  int u = stk.back(); stk.pop_back();
                                                                  int v = stk.back(); stk.pop_back();
  if(dis[edd]==INF_WEI) return {-1, -1};
                                                                  match[u] = v;
 Cap mw=INF_CAP;
                                                                  match[v] = u;
 for(int i=edd;i!=ori;i=fa[i])
  mw=min(mw,G[fa[i]][wh[i]].cap);
 for (int i=edd;i!=ori;i=fa[i]){
  auto &eg=G[fa[i]][wh[i]];
                                                               if (!found) break;
```

```
int ret = 0;
for (int i=0; i<n; i++)
  ret += edge[i][match[i]];
  return ret>>1;
}
graph;
```

# 5 Math

# 5.1 $\lfloor \frac{n}{i} \rfloor$ Enumeration

$$T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T \cdot + 1} \rfloor} \rfloor$$

# 5.2 Strling Number

#### 5.2.1 First Kind

 $S_1(n,k)$  counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1) \dots (x+n-1) = \sum_{k=0}^n S_1(n,k) x^k$$

$$g(x) = x(x+1) \dots (x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

#### 5.2.2 Second Kind

 $S_2(n,k)$  counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k \binom{k}{i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

# 5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b==(x/y)*a;
}
```

# 5.4 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
  vector<T> d(output.size() + 1), me, he;
  for (size_t f = 0, i = 1; i <= output.size(); ++i) {
    for (size_t j = 0; j < me.size(); ++j)
        d[i] += output[i - j - 2] * me[j];
    if ((d[i] -= output[i - 1]) == 0) continue;
    if (me.empty()) {
        me.resize(f = i);
        continue;
    }
    vector<T> o(i - f - 1);
    T k = -d[i] / d[f]; o.push_back(-k);
    for (T x : he) o.push_back(x * k);
    if (o.size() < me.size()) o.resize(me.size());
    for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];
    if (i-f+he.size() >= me.size()) he = me, f = i;
    me = o;
    }
    return me;
}
```

### 5.5 Charateristic Polynomial

```
vector<vector<int>>> &A) {
  int N = A.size();
  vector<vector<int>>> h = A;
  for (int i = 0; i < N - 2; ++i) {
    if (!H[i + 1][i]) {
      for (int j = i + 2; j < N; ++j) {
        if (H[j][i]) {
            for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j][k]);
        }
}</pre>
```

```
for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k]
    ][j]);
     break:
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;
  for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL</pre>
     * H[i + 1][k] * (kP - coef)) % kP;
   for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i + 1] + 1LL * H[k][j] * coef) % kP;
 return H;
vector<int> CharacteristicPoly(const vector<vector<int
    >> &A) {
 int N = A.size();
 auto H = Hessenberg(A);
 for (int i = 0; i < N; ++i) {
  for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
 vector<vector<int>> P(N + 1, vector<int>(N + 1));
 P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
  P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
    1];
  int val = 1;
  for (int j = i - 1; j >= 0; --j) {
   int coef = 1LL * val * H[j][i - 1] % kP;
   for (int k = 0; k \le j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 if (N & 1) {
  for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
 return P[N];
5.6 Chinese Remainder
1ld crt(lld ans[], lld pri[], int n){
 11d M = 1, ret = 0;
 for(int i=0;i<n;i++) M *= pri[i];
 for(int i=0;i<n;i++){</pre>
  lld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
  ret += (ans[i]*(M/pri[i])%M * iv)%M;
  ret %= M;
 return ret;
/*
Another:
x = a1 \% m1
x = a2 \% m2
g = gcd(m1, m2)
assert((a1-a2)\%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
0 <= x < lcm(m1, m2)
5.7 De-Bruijn
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
 if (t > n) {
  if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
 } else {
  aux[t] = aux[t - p];
  db(t + 1, p, n, k);
for (int i = aux[t - p] + 1; i < k; ++i) {
```

aux[t] = i;

db(t + 1, t, n, k);

```
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int de_bruijn(int k, int n) {
 // return cyclic string of len k^n s.t. every string
  // of len n using k char appears as a substring.
 if (k == 1) {
   res[0] = 0;
   return 1:
 for (int i = 0; i < k * n; i++) aux[i] = 0;
 sz = 0;
 db(1, 1, n, k);
 return sz;
5.8 DiscreteLog
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
   // x^? \equiv y (mod M)
   Int t = 1, c = 0, g = 1;
   for (Int M<sub>_</sub> = M; M<sub>_</sub> > 0; M<sub>_</sub> >>= 1)
      g = g * x % M;
   for (g = gcd(g, M); t % g != 0; ++c) {
      if (t == y) return c;
      t = t * x % M;
   if (y % g != 0) return -1;
  t /= g, y /= g, M /= g;

Int h = 0, gs = 1;

for (; h * h < M; ++h) gs = gs * x % M;
   unordered_map<Int, Int> bs;
   for (Int s = 0; s < h; bs[y] = ++s)
     y = y * x % M;
   for (Int s = 0; s < M; s += h) {
      t = t * gs % M;
      if (bs.count(t)) return c + s + h - bs[t];
   return -1;
}
5.9 Extended Euler
       a^b \equiv \begin{cases} a^b \mod \varphi(m) + \varphi(m) & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases}
                                                                         (\mathsf{mod}\ m)
5.10 ExtendedFloorSum
  g(a,b,c,n) = \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor
                     \begin{array}{l} \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), \end{array}
                                                                         a \geq c \vee b \geq c
                                                                         n < 0 \lor a = 0
                     \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
                    (-h(c, c-b-1, a, m-1)),
                                                                         otherwise
  h(a, b, c, n) = \sum_{i=0}^{n} \left\lfloor \frac{ai + b}{c} \right\rfloor^{2}
                    \left( \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right)
                     +\lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot \tilde{n}(n+1)
                     +h(a \bmod c, b \bmod c, c, n)
```

# 5.11 Fast Fourier Transform

```
const int mod = 1000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
    static_assert (M1 <= M2 && M2 <= M3);
    constexpr int64_t r12 = modpow(M1, M2-2, M2);
    constexpr int64_t r13 = modpow(M1, M3-2, M3);
    constexpr int64_t r23 = modpow(M2, M3-2, M3);
    constexpr int64_t M1M2 = 1LL * M1 * M2 * mod;
    B = (B - A + M2) * r12 * M2;
    C = (C - A + M3) * r13 * M3;
    C = (C - B + M3) * r23 * M3;
    return (A + B * M1 + C * M1M2) * mod;
}</pre>
```

 $+2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)$ 

 $+2\lfloor \frac{\bar{b}}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),$ 

nm(m+1) - 2g(c, c-b-1, a, m-1)

-2f(c,c-b-1,a,m-1)-f(a,b,c,n), otherwise

 $\begin{aligned} a &\geq c \vee b \geq c \\ n &< 0 \vee a = 0 \end{aligned}$ 

```
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
for (int i = 0; i <= maxn; i++)</pre>
  omega[i] = cplx(cos(2 * pi * j / maxn),
      sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
int z = __builtin_ctz(n) - 1;
for (int i = 0; i < n; ++i) {</pre>
  int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^{=(i >> j & 1)<<(z - j);
  if (x > i) swap(v[x], v[i]);
 for (int s = 2; s <= n; s <<= 1) {
  int z = s \gg 1;
  for (int i = 0; i < n; i += s) {
   for (int k = 0; k < z; ++k) {
  cplx x = v[i + z + k] * omega[maxn / s * k];
    v[i + z + k] = v[i + k] - x;
    v[i+k] = v[i+k] + x;
void ifft(vector<cplx> &v, int n) {
 fft(v, n); reverse(v.begin() + 1, v.end());
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
VL convolution(const VI &a, const VI &b) {
 // Should be able to handle N <= 10^5, C <= 10^4
 int sz = 1;
 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 vector<cplx> v(sz);
 for (int i = 0; i < sz; ++i) {
  double re = i < a.size() ? a[i] : 0;</pre>
  double im = i < b.size() ? b[i] : 0;</pre>
  v[i] = cplx(re, im);
 fft(v, sz);
 for (int i = 0; i <= sz / 2; ++i) {
  int j = (sz - i) & (sz - 1);</pre>
  cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
  * cplx(0, -0.25);
if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
     ].conj()) * cplx(0, -0.25);
  v[i] = x;
 ifft(v, sz);
 VL c(sz);
 for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
 return c:
VI convolution_mod(const VI &a, const VI &b, int p) {
 int sz = 1;
 while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
 vector<cplx> fa(sz), fb(sz);
for (int i = 0; i < (int)a.size(); ++i)</pre>
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 for (int i = 0; i < (int)b.size(); ++i)
fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
 fft(fa, sz), fft(fb, sz);
 double r = 0.25 / sz;
 cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
  int j = (sz - i) & (sz - 1);
  cplx a1 = (fa[i] + fa[j].conj());
  cplx a2 = (fa[i] - fa[j].conj()) * r2;
  cplx b1 = (fb[i] + fb[j].conj()) * r3;
  cplx b2 = (fb[i] - fb[j].conj()) * r4;
  if (i != j) {
   cplx c1 = (fa[j] + fa[i].conj());
cplx c2 = (fa[j] - fa[i].conj()) * r2;
cplx d1 = (fb[j] + fb[i].conj()) * r3;
   cplx d2 = (fb[j] - fb[i].conj()) * r4;
   fa[i] = c1 * d1 + c2 * d2 * r5;
   fb[i] = c1 * d2 + c2 * d1;
```

```
5.14 Gauss Elimination
  fa[j] = a1 * b1 + a2 * b2 * r5;
  fb[j] = a1 * b2 + a2 * b1;
                                                              void gauss(vector<vector<double>> &d) {
                                                               int n = d.size(), m = d[0].size();
 fft(fa, sz), fft(fb, sz);
                                                               for (int i = 0; i < m; ++i) {
 vector<int> res(sz);
                                                                int p = -1;
 for (int i = 0; i < sz; ++i) {
                                                                for (int j = i; j < n; ++j) {
  long long a = round(fa[i].re), b = round(fb[i].re),
                                                                 if (fabs(d[j][i]) < eps) continue;</pre>
       c = round(fa[i].im);
                                                                 if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
  res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
                                                                if (p == -1) continue;
 return res;
                                                                for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
for (int j = 0; j < n; ++j) {
  if (i == j) continue;</pre>
}}
5.12 FloorSum
                                                                 double z = d[j][i] / d[i][i];
                                                                 for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
// @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
 11u ans = 0;
                                                                     Miller Rabin
                                                              5.15
 while (true)
 if (a >= m) {
                                                              bool isprime(llu x){
   ans += n * (n - 1) / 2 * (a / m); a %= m;
                                                               static llu magic[]={2,325,9375,28178,\
                                                                         450775,9780504,1795265022};
  if (b >= m) {
                                                               static auto witn=[](llu a,llu u,llu n,int t)
   ans += n * (b / m); b %= m;
                                                                if (!(a = mpow(a%n,u,n)))return 0;
  llu y_max = a * n + b;
                                                                while(t--)
  if (y_max < m) break;</pre>
                                                                 1lu a2=mul(a,a,n);
  // y_max < m * (n + 1)
                                                                 if(a2==1 && a!=1 && a!=n-1)
 // floor(y_max / m) <= n
                                                                   return 1:
 n = (11u)(y_max / m), b = (11u)(y_max % m);
                                                                 a = a2;
  swap(m, a);
                                                                }
                                                                return a!=1;
 return ans;
                                                               if(x<2)return 0;</pre>
11d floor_sum(1ld n, 1ld m, 1ld a, 1ld b) {
                                                                if(!(x&1))return x==2;
 llu ans = 0;
                                                               11u x1=x-1:int t=0:
 if (a < 0) {
                                                               while(!(x1&1))x1>>=1,t++;
 llu a2 = (a % m + m) % m;
                                                               for(llu m:magic)if(witn(m,x1,x,t))return 0;
  ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
                                                               return 1;
 if (b < 0) {
                                                              5.16 NTT
 11u b2 = (b \% m + m) \% m;
                                                              template <int mod, int G, int maxn>
  ans -= 1ULL * n * ((b2 - b) / m);
                                                              struct NTT {
                                                               static_assert (maxn == (maxn & -maxn));
                                                               int roots[maxn];
 return ans + floor_sum_unsigned(n, m, a, b);
                                                               NTT () {
                                                                int r = modpow(G, (mod - 1) / maxn);
                                                                for (int i = maxn >> 1; i; i >>= 1) {
5.13 FWT
                                                                 roots[i] = 1;
                                                                 for (int j = 1; j < i; j++)
roots[i + j] = modmul(roots[i + j - 1], r);</pre>
/* xor convolution:
 * x = (x0,x1) , y = (y0,y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                  r = modmul(r, r);
                                                                }
 * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
 * z' = ( ( x0+x1 )( y0+y1 ) , ( x0-x1 )( y0-y1 ) )
* z = (1/2) * z''
                                                               // n must be 2^k, and 0 \le F[i] < mod
                                                               void operator()(int F[], int n, bool inv = false) {
 * or convolution:
                                                                for (int i = 0, j = 0; i < n; i++) {
  if (i < j) swap(F[i], F[j]);
  for (int b)</pre>
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                                 for (int k = n > 1; (j^k < k) < k; k > = 1);
 * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
const LL MOD = 1e9+7;
                                                                for (int s = 1; s < n; s *= 2) {
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
                                                                 for (int i = 0; i < n; i += s * 2) {
 for( int d = 1 ; d < N ; d <<= 1 ) {
                                                                   for (int j = 0; j < s; j++) {
  int d2 = d << 1;
                                                                    int a = F[i+j];
  for( int s = 0 ; s < N ; s += d2 )
                                                                    int b = modmul(F[i+j+s], roots[s+j]);
   for( int i = s , j = s+d ; i < s+d ; i++, j++ ){</pre>
                                                                    F[i+j] = modadd(a, b); // a + b
   LL ta = x[ i ] , tb = x[ j ];
                                                                    F[i+j+s] = modsub(a, b); // a - b
    x[i] = ta+tb;
    x[ j ] = ta-tb;
                                                                 }
    if( x[ i ] >= MOD ) x[ i ] -= MOD;
    if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
                                                                if (inv) {
                                                                 int invn = modinv(n);
                                                                  for (int i = 0; i < n; i++)
 if( inv )
                                                                   F[i] = modmul(F[i], invn);
  for( int i = 0 ; i < N ; i++ ) {</pre>
                                                                 reverse(F + 1, F + n);
   x[i] *= inv(N, MOD);
   x[ i ] %= MOD;
                                                              NTT<2013265921, 31, 1048576> ntt;
```

```
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5.17 Range Sieve
const int MAX_SQRT_B = 50000;
const int MAX_L = 200000 + 5;
bool is_prime_small[MAX_SQRT_B], is_prime[MAX_L];
void sieve(lld l, lld r)\{ // [l, r)
 for(lld i=2;i*i<r;i++) is_prime_small[i] = true;</pre>
 for(lld i=l;i<r;i++) is_prime[i-l] = true;</pre>
 if(l==1) is_prime[0] = false;
 for(lld i=2;i*i<r;i++){</pre>
  if(!is_prime_small[i]) continue;
  for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;</pre>
  for(1ld j=std::max(2LL, (1+i-1)/i)*i;j<r;j+=i)</pre>
   is_prime[j-l]=false;
}
      Partition Number
5.18
int b = sqrt(n);
ans[0] = tmp[0] = 1;
for (int i = 1; i <= b; i++) {
for (int rep = 0; rep < 2; rep++)</pre>
 for (int j = i; j <= n - i * i; j++)
modadd(tmp[j], tmp[j-i]);
for (int j = i * i; j <= n; j++)</pre>
  modadd(ans[j], tmp[j - i * i]);
5.19 Pi Count (Linear Sieve)
static constexpr int N = 1000000 + 5;
11d pi[N];
vector<int> primes;
bool sieved[N];
11d cube_root(11d x){
 1ld s=cbrt(x-static_cast<long double>(0.1));
 while(s*s*s <= x) ++s;
 return s-1;
11d square_root(11d x){
 lld s=sqrt(x-static_cast<long double>(0.1));
 while(s*s <= x) ++s;</pre>
 return s-1;
void init(){
 primes.reserve(N);
 primes.push_back(1);
 for(int i=2;i<N;i++) {</pre>
  if(!sieved[i]) primes.push_back(i);
  pi[i] = !sieved[i] + pi[i-1];
  for(int p: primes) if(p > 1) {
  if(p * i >= N) break;
   sieved[p * i] = true;
   if(p % i == 0) break;
11d phi(11d m, 11d n) {
 static constexpr int MM = 80000, NN = 500;
 static 1ld val[MM][NN];
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
 1ld ret = phi(m,n-1)-phi(m/primes[n],n-1);
 if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
 return ret;
1ld pi_count(lld);
11d P2(11d m, 11d n) {
 1ld sm = square_root(m), ret = 0;
 for(lld i = n+1;primes[i]<=sm;i++)</pre>
  ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
 return ret;
11d pi_count(11d m) {
 if(m < N) return pi[m];</pre>
 11d n = pi_count(cube_root(m));
 return phi(m, n) + n - 1 - P2(m, n);
5.20 Pollard Rho
// does not work when n is prime
```

// return any non-trivial factor

```
llu pollard_rho(llu n){
    static auto f=[](llu x,llu k,llu m){
        return add(k,mul(x,x,m),m);};
    if (!(n&1)) return 2;
    mt19937 rnd(120821011);
    while(true){
        llu y=2,yy=y,x=rnd()%n,t=1;
        for(llu sz=2;t==1;sz<<=1,y=yy) {
            for(llu i=0;t==1&&i<sz;++i){
                 yy=f(yy,x,n);
                 t=gcd(yy>y?yy-y:y-yy,n);
            }
        if(t!=1&&t!=n) return t;
    }
}
```

# 5.21 Polynomial Operations

```
using V = vector<int>;
#define fi(s, n) for (int i = int(s); i < int(n); ++i)
template <int mod, int G, int maxn> struct Poly : V {
 static uint32_t n2k(uint32_t n) {
  if (n <= 1) return 1;
  return 1u << (32 - __builtin_clz(n - 1));</pre>
 static NTT<mod,G,maxn> ntt; // coefficients in [0, P)
 explicit Poly(int n = 1) : V(n) {}
 Poly(const V &v) : V(v) {}
 Poly(const Poly &p, size_t n) : V(n) {
  copy_n(p.data(), min(p.size(), n), data());
 Poly &irev() { return reverse(data(), data() + size())
     , *this; }
 Poly &isz(int _n) { return resize(_n), *this; }
 Poly &iadd(const Poly &rhs) { // n() == rhs.n()
  fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
  return *this;
 Poly &imul(int k) {
  fi(0, size())(*this)[i] = modmul((*this)[i], k);
  return *this;
 Poly Mul(const Poly &rhs) const {
  const int _n = n2k(size() + rhs.size() - 1);
  Poly X(*this, _n), Y(rhs, _n);
ntt(X.data(), _n), ntt(Y.data(),
  fi(0, _n) X[i] = modmul(X[i], Y[i]);
  ntt(X.data(), _n, true);
  return X.isz(size() + rhs.size() - 1);
 Poly Inv() const { // coef[0] != 0
  if (size() == 1) return V{modinv(*begin())};
  const int _n = n2k(size() * 2);
  Poly Xi = Poly(*this, (size() + 1) / 2).Inv().isz(_n)
      Y(*this, _n);
  ntt(Xi.data(), _n), ntt(Y.data(), _n);
fi(0, _n) Xi[i] = modmul(Xi[i], modsub(2, modmul(Xi[i], Y[i])));
  ntt(Xi.data(), _n, true);
  return Xi.isz(size());
 Poly Sqrt() const { // coef[0] \in [1, mod)^2
  if (size() == 1) return V{QuadraticResidue((*this)
    [0], mod)};
  Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
    size())
  return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
    + 1);
 pair<Poly, Poly> DivMod(const Poly &rhs) const {
  if (size() < rhs.size()) return {V{0}, *this};</pre>
  const int _n = size() - rhs.size() + 1;
  Poly X(rhs); X.irev().isz(_n);
  Poly Y(*this); Y.irev().isz(_n);
  Poly Q = Y.Mul(X.Inv()).isz(_n).irev();
  X = rhs.Mul(Q), Y = *this;
fi(0, size()) Y[i] = modsub(Y[i], X[i]);
  return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
 Poly Dx() const {
  Poly ret(size() - 1);
```

```
fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
                                                                  Poly 0 = R.Mul(Poly(data(), min<int>(2 << level,
                                                                   size())));
   1]);
                                                                  Poly Q(2 << level); Q[0] = 1;
 return ret.isz(max<int>(1, ret.size()));
                                                                  for (int j = (1 << level); j < (2 << level); ++j)</pre>
                                                                   Q[j] = modsub(mod, O[j]);
Poly Sx() const {
 Poly ret(size() + 1);
                                                                  R = R.Mul(Q).isz(4 << level);
 fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
   this)[i]);
                                                                 return R.isz(L);
 return ret;
                                                               static int LinearRecursion(const V &a, const V &c,
Poly _tmul(int nn, const Poly &rhs) const {
                                                                   int64_t n) { // a_n = \sum_{i=1}^{n} a_{i}(n-j)}
 Poly Y = Mul(rhs).isz(size() + nn - 1);
                                                                 const int k = (int)a.size();
 return V(Y.data() + size() - 1, Y.data() + Y.size());
                                                                assert((int)c.size() == k + 1);
Poly C(k + 1), W({1}, k), M = {0, 1};
                                                                 fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
V _eval(const V &x, const vector<Poly> &up) const {
                                                                 C[k] = 1;
 const int _n = (int)x.size();
 if (!_n) return {};
                                                                 while (n) {
                                                                  if (n % 2) W = W.Mul(M).DivMod(C).second;
 vector<Poly> down(_n * 2);
 // down[1] = DivMod(up[1]).second;
                                                                 n /= 2, M = M.Mul(M).DivMod(C).second;
 // fi(2, _n * 2) down[i] = down[i / 2].DivMod(up[i]).
                                                                 int ret = 0:
   second:
 down[1] = Poly(up[1]).irev().isz(size()).Inv().irev()
                                                                 fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
 ._tmul(_n, *this);
fi(2, _n * 2) down[i] = up[i ^ 1]._tmul(up[i].size()
                                                                 return ret;
   - 1, down[i / 2]);
 V y(_n); fi(0, _n) y[i] = down[_n + i][0];
                                                              #undef fi
                                                              using Poly_t = Poly<998244353, 3, 1 << 20>;
 return y;
                                                              template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
static vector<Poly> _tree1(const V &x) {
                                                              5.22 Quadratic residue
 const int _n = (int)x.size();
 vector<Poly> up(_n * 2);
                                                              struct S {
                                                               int MOD, w;
 fi(0, _n) up[_n + i] = V\{modsub(mod, x[i]), 1\};
 for(int i=_n-1;i>0;--i) up[i] = up[i * 2].Mul(up[i *
                                                               int64_t x, y;
   2 + 1]);
                                                               S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
 return up;
                                                                 : MOD(m), w(w_{-}), x(x_{-}), y(y_{-}) {}
                                                                S operator*(const S &rhs) const {
V Eval(const V &x) const { return _eval(x, _tree1(x));
                                                                 int w_ = w;
                                                                 if (w<sub>_</sub> == -1) w<sub>_</sub> = rhs.w;
static Poly Interpolate(const V &x, const V &y) {
                                                                 assert(w_! = -1 \text{ and } w_! = rhs.w);
                                                                 return \{ MOD, w_,
 const int _n = (int)x.size();
 vector<Poly> up = _tree1(x), down(_n * 2);
                                                                  (x * rhs.x + y * rhs.y % MOD * w) % MOD,
 V z = up[1].Dx()._eval(x, up);
                                                                  (x * rhs.y + y * rhs.x) % MOD };
 fi(0, _n) z[i] = modmul(y[i], modinv(z[i]));
 fi(0, _n) down[_n + i] = V{z[i]};
 for(int i=_n-1;i>0;--i) down[i] = down[i * 2].Mul(up[
    i * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
                                                              int get_root(int n, int P) {
                                                                if (P == 2 or n == 0) return n;
 return down[1];
                                                                 if (qpow(n, (P - 1) / 2, P) != 1) return -1;
                                                                 auto check = [&](int x) {
Poly Ln() const \{ // coef[0] == 1 \}
                                                                 return qpow(x, (P - 1) / 2, P); };
if (check(n) == P-1) return -1;
return Dx().Mul(Inv()).Sx().isz(size());
                                                                 int64_t a; int w; mt19937 rnd(7122);
Poly Exp() const { // coef[0] == 0
                                                                 do { a = rnd() % P;
                                                                   w = ((a * a - n) % P + P) % P;
 if (size() == 1) return V{1};
 Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size
                                                                 } while (check(w) != P - 1);
                                                                 return qpow(S(P, w, a, 1), (P + 1) / 2).x;
   ());
 Poly Y = X.Ln(); Y[0] = mod - 1;
 fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
                                                              5.23 Simplex
 return X.Mul(Y).isz(size());
                                                              namespace simplex {
Poly Pow(const string &K) const {
                                                              // maximize c^Tx under Ax <= B
                                                              // return VD(n, -inf) if the solution doesn't exist // return VD(n, +inf) if the solution is unbounded
 int nz = 0:
 while (nz < size() && !(*this)[nz]) ++nz;</pre>
 int nk = 0, nk2 = 0;
                                                              using VD = vector<double>;
 for (char c : K) {
                                                              using VVD = vector<vector<double>>;
  nk = (nk * 10 + c - '0') % mod;
                                                              const double eps = 1e-9;
  nk2 = nk2 * 10 + c - '0';
                                                              const double inf = 1e+9;
  if (nk2 * nz >= size())
                                                              int n, m;
   return Poly(size());
                                                              VVD d;
                                                              vector<int> p, q;
  nk2 \% = mod - 1;
                                                              void pivot(int r, int s) {
                                                               double inv = 1.0 / d[r][s];
 if (!nk && !nk2) return Poly(V{1}, size());
 Poly X = V(data() + nz, data() + size() - nz * (nk2 -
                                                               for (int i = 0; i < m + 2; ++i)
                                                                for (int j = 0; j < n + 2; ++j)
if (i != r && j != s)
    1));
 int x0 = X[0];
 return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
                                                                   d[i][j] = d[r][j] * d[i][s] * inv;
                                                               for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
   modpow(x0, nk2)).irev().isz(size()).irev();
Poly InvMod(int L) { // (to evaluate linear recursion)
                                                               d[r][s] = inv; swap(p[r], q[s]);
 Poly R{1, 0}; // *this * R mod x^L = 1 (*this[0] ==
                                                              bool phase(int z) {
 for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                               int x = m + z;
                                                               while (true) {
```

```
int s = -1;
  for (int i = 0; i <= n; ++i) {
   if (!z && q[i] == -1) continue;
   if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
  if (d[x][s] > -eps) return true;
  int r = -1;
  for (int i = 0; i < m; ++i) {
   if (d[i][s] < eps) continue;</pre>
   if (r == -1 ||
    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
  pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
m = b.size(), n = c.size();
 d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)</pre>
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)
p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
 for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
 q[n] = -1, d[m + 1][n] = 1;
 int r = 0:
 for (int i = 1; i < m; ++i)</pre>
  if (d[i][n + 1] < d[r][n + 1]) r = i;
 if (d[r][n + 1] < -eps) {</pre>
  pivot(r, n);
  if (!phase(1) || d[m + 1][n + 1] < -eps)
    return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
   int s = min_element(d[i].begin(), d[i].end() - 1)
         - d[i].begin();
   pivot(i, s);
  }
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)
if (p[i] < n) x[p[i]] = d[i][n + 1];
 return x;
}}
5.24 Simplex Construction
\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j. \text{and } x_i \geq 0 \text{ for all } 1 \leq i \leq n.
  1. In case of minimization, let c'_i = -c_i
  2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
```

Standard form: maximize  $\sum_{1 \leq i \leq n} c_i x_i$  such that for all  $1 \leq j \leq m$ ,

- $3. \sum_{1 \le i \le n} A_{ji} x_i = b_j$ 
  - $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$
  - $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

# Geometry

#### 6.1 Basic Geometry

```
#define IM imag
#define RE real
using 1ld = int64_t;
using llf = long double;
using Point = std::complex<lld>;
using Pointf = std::complex<11f>;
auto toPointf(Point p) { return Pointf{IM(p), RE(p)}; }
int sgn(11d x) \{ return (x > 0) - (x < 0); \}
1ld dot(Point a, Point b) { return RE(conj(a) * b); }
lld cross(Point a, Point b) { return IM(conj(a) * b); }
int ori(Point a, Point b, Point c) {
return sgn(cross(b - a, c - a));
bool operator<(const Point &a, const Point &b) {</pre>
return RE(a) != RE(b) ? RE(a) < RE(b) : IM(a) < IM(b);
int argCmp(Point a, Point b) {
```

```
// -1 / 0 / 1 <-> < / == / > (atan2)
 int qa = (IM(a) == 0
   ? (RE(a) < 0 ? 3 : 1) : (IM(a) < 0 ? 0 : 2));
 int qb = (IM(b) == 0
   ? (RE(b) < 0 ? 3 : 1) : (IM(b) < 0 ? 0 : 2));
 if (qa != qb)
  return sgn(qa - qb);
 return sgn(cross(b, a));
template <typename V> llf area(const V & pt) {
 11d ret = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 return ret / 2.0;
Point rot90(Point p) { return Point{-IM(p), RE(p)}; }
Pointf project(Pointf p, Pointf q) { // p onto q return dot(p, q) * q / dot(q, q);
```

# 6.2 Segment & Line Intersection

```
struct Segment {
 Point st, dir; // represent st + t*dir for 0<=t<=1
 Segment(Point s, Point e) : st(s), dir(e - s) {}
 static bool valid(lld p, lld q)
  // is there t s.t. 0 <= t <= 1 && qt == p ?
  if (q < 0) q = -q, p = -p;
  return 0 <= p && p <= q;
};
bool isInter(Segment A, Point P) {
if (A.dir == Point(0)) return P == A.st;
 return cross(P - A.st, A.dir) == 0 &&
  Segment::valid(dot(P - A.st, A.dir), norm(A.dir));
template <typename U, typename V>
bool isInter(U A, V B) {
  if (cross(A.dir, B.dir) == 0)
  return // handle parallel yourself
   isInter(A, B.st) || isInter(A, B.st+B.dir) ||
   isInter(B, A.st) || isInter(B, A.st+A.dir);
 Point D = B.st - A.st;
 11d C = cross(A.dir, B.dir)
 return U::valid(cross(D, A.dir), C) &&
   V::valid(cross(D, B.dir), C);
struct Line {
Point st, ed, dir;
 Line (Point s, Point e)
  : st(s), ed(e), dir(e - s) {}
Pointf intersect(const Line &A, const Line &B) {
11f t = cross(B.st - A.st, B.dir) /
 llf(cross(A.dir, B.dir))
 return toPointf(A.st) +
  Pointf(t) * toPointf(A.dir);
```

#### 2D Convex Hull 6.3

```
template<typename PT>
vector<PT> buildConvexHull(vector<PT> d) {
 sort(ALL(d), [](const PT& a, const PT& b){
   return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
 vector<PT> s(SZ(d)<<1);</pre>
 int o = 0;
 for(auto p: d) {
  while(o \ge 2 \& cross(p-s[o-2], s[o-1]-s[o-2]) <= 0)
  s[o++] = p;
 for(int i=SZ(d)-2, t = o+1; i>=0; i--)
  while(o>=t\&cross(d[i]-s[o-2],s[o-1]-s[o-2])<=0)
   0--
  s[o++] = d[i];
 s.resize(o-1);
 return s;
```

## 6.4 3D Convex Hull

```
6.7 kD Closest Pair (3D ver.)
// return the faces with pt indexes
int flag[MXN][MXN];
                                                                11f solve(vector<P> v) {
struct Point{
                                                                 shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,</pre>
 ld x,y,z;
 Point operator * (const ld &b) const {
                                                                  unordered_map<lld, int>>> m;
  return (Point){x*b,y*b,z*b};}
                                                                 llf d = dis(v[0], v[1]);
 Point operator * (const Point &b) const {
                                                                 auto Idx = [&d] (11f x) -> 11d {
  return(Point){y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
                                                                  return round(x * 2 / d) + 0.1; };
                                                                 auto rebuild_m = [&m, &v, &Idx](int k) {
                                                                  m.clear();
Point ver(Point a, Point b, Point c) {
                                                                  for (int i = 0; i < k; ++i)
return (b - a) * (c - a);}
                                                                   m[Idx(v[i].x)][Idx(v[i].y)]
vector<Face> convex_hull_3D(const vector<Point> pt) {
                                                                    [Idx(v[i].z)] = i;
 int n = SZ(pt), ftop = 0
                                                                 }; rebuild_m(2);
 REP(i,n) REP(j,n) flag[i][j] = 0;
                                                                 for (size_t i = 2; i < v.size(); ++i) {
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),</pre>
 vector<Face> now;
 now.emplace_back(0,1,2);
                                                                     kz = Idx(v[i].z); bool found = false;
 now.emplace_back(2,1,0);
                                                                  for (int dx = -2; dx <= 2; ++dx) {
 for (int i=3; i<n; i++){
                                                                   const 11d nx = dx + kx;
  ftop++; vector<Face> next;
                                                                   if (m.find(nx) == m.end()) continue;
  REP(j, SZ(now)) {
  Face& f=now[j]; int ff = 0;
                                                                   auto\& mm = m[nx];
                                                                   for (int dy = -2; dy <= 2; ++dy) {
   ld d=(pt[i]-pt[f.a]).dot(
                                                                    const 11d ny = dy + ky;
     ver(pt[f.a], pt[f.b], pt[f.c]));
                                                                    if (mm.find(ny) == mm.end()) continue;
   if (d <= 0) next.push_back(f);</pre>
                                                                    auto& mmm = mm[ny];
   if (d > 0) ff=ftop;
else if (d < 0) ff=-ftop;</pre>
                                                                    for (int dz = -2; dz <= 2; ++dz) {
                                                                     const lld nz = dz + kz;
   flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
                                                                     if (mmm.find(nz) == mmm.end()) continue;
                                                                     const int p = mmm[nz];
  REP(j, SZ(now)) {
                                                                     if (dis(v[p], v[i]) < d) {</pre>
   Face& f=now[j];
                                                                      d = dis(v[p], v[i]);
   if (flag[f.a][f.b] > 0 &&
                                                                      found = true;
     flag[f.a][f.b] != flag[f.b][f.a])
    next.emplace_back(f.a,f.b,i);
                                                                    }
   if (flag[f.b][f.c] > 0 &&
     flag[f.b][f.c] != flag[f.c][f.b])
    next.emplace_back(f.b,f.c,i);
                                                                  if (found) rebuild_m(i + 1);
   if (flag[f.c][f.a] > 0 &&
                                                                  else m[kx][ky][kz] = i;
     flag[f.c][f.a] != flag[f.a][f.c])
    next.emplace_back(f.c,f.a,i);
                                                                 return d;
                                                               }
 now=next;
                                                               6.8 Simulated Annealing
 return now;
                                                               11f anneal() {
                                                                mt19937 rnd_engine( seed );
                                                                 uniform_real_distribution< llf > rnd( 0, 1 );
      2D Farthest Pair
                                                                 const llf dT = 0.001;
                                                                 // Argument p
// stk is from convex hull
                                                                 11f S_cur = calc( p ), S_best = S_cur;
for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
                                                                  // Modify p to p_prime
const llf S_prime = calc( p_prime );
for(int i=0;i<n;i++) {</pre>
 while(abs(cross(stk[i+1]-stk[i],
                                                                  const llf delta_c = S_prime - S_cur
   stk[(pos+1)%n]-stk[i])) >
                                                                  11f prob = min( ( 11f ) 1, exp( -delta_c / T ) );
   abs(cross(stk[i+1]-stk[i]
                                                                  if ( rnd( rnd_engine ) <= prob )</pre>
   stk[pos]-stk[i]))) pos = (pos+1)%n;
                                                                   S_cur = S_prime, p = p_prime;
 ans = max({ans, dis(stk[i], stk[pos]),
                                                                  if ( S_prime < S_best ) // find min</pre>
  dis(stk[i+1], stk[pos])});
                                                                   S_best = S_prime, p_best = p_prime;
                                                                 return S_best;
6.6 2D Closest Pair
                                                               }
struct cmp_y {
                                                                6.9 Half Plane Intersection
 bool operator()(const P& p, const P& q) const {
  return p.y < q.y;</pre>
                                                                // cross(pt-line.st, line.dir)<=0 <-> pt in half plane
                                                               bool operator<(const Line &lhs, const Line &rhs) {</pre>
                                                                  if (int cmp = argCmp(lhs.dir, rhs.dir))
multiset<P, cmp_y> s;
                                                                    return cmp == -1;
void solve(P a[], int n) {
                                                                  return ori(lhs.st, lhs.ed, rhs.st) < 0;
 sort(a, a + n, [](const P& p, const P& q) {
                                                               }
  return tie(p.x, p.y) < tie(q.x, q.y);</pre>
                                                                // intersect function is in "Segment Intersect"
 11f d = INF; int pt = 0;
for (int i = 0; i < n; ++i) {</pre>
                                                               llf HPI(vector<Line> &lines) {
                                                                  sort(lines.begin(), lines.end());
  while (pt < i \text{ and } a[i].x - a[pt].x >= d)
                                                                  deque<Line> que;
   s.erase(s.find(a[pt++]))
                                                                  deque<Pointf> pt;
  auto it = s.lower_bound(P(a[i].x, a[i].y - d));
                                                                  que.push_back(lines[0]);
for (int i = 1; i < (int)lines.size(); i++) {</pre>
  while (it != s.end() and it->y - a[i].y < d)</pre>
   d = min(d, dis(*(it++), a[i]));
                                                                    if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
  s.insert(a[i]);
```

continue; #define POP(L, R) \

while (pt.size() > 0 \

double S, h, theta;

```
&& ori(L.st, L.ed, pt.back()) < 0) \
                                                               double a=abs(pb),b=abs(pa),c=abs(pb-pa);
                                                               double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
      pt.pop_back(), que.pop_back(); \
    while (pt.size() > 0 \
                                                               double cosC = dot(pa,pb) / a / b, C = acos(cosC);
      && ori(R.st, R.ed, pt.front()) < 0) \
                                                               if(a > r){
      pt.pop_front(), que.pop_front();
                                                                S = (C/2)*r*r;
                                                                h = a*b*sin(C)/c;
    POP(lines[i], lines[i]);
    pt.push_back(intersect(que.back(), lines[i]));
                                                                if (h < r && B < PI/2)
                                                                 S = (acos(h/r)*r*r - h*sqrt(r*r-h*h));
    que.push_back(lines[i]);
 POP(que.front(), que.back())
                                                               else if(b > r){
                                                                theta = PI - B - asin(sin(B)/r*a);
  if (que.size() <= 1 ||</pre>
    argCmp(que.front().dir, que.back().dir) == 0)
                                                                S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
    return 0:
  pt.push_back(intersect(que.front(), que.back()));
                                                               else S = .5*sin(C)*a*b;
                                                               return S;
  return area(pt);
                                                              double area_poly_circle(const vector<pdd> &poly,
     Minkowski Sum
6.10
                                                                const pdd &0.const double r){
                                                               double S=0; int N=poly.size();
vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
                                                               for(int i=0;i<N;++i)</pre>
hull(A), hull(B);
                                                                S += _area(poly[i]-0, poly[(i+1)%N]-0, r)
 vector<pll> C(1, A[0] + B[0]), s1, s2;
                                                                  * ori(0, poly[i], poly[(i+1)%N]);
for(int i = 0; i < SZ(A); ++i)
                                                               return fabs(S);
 s1.pb(A[(i + 1) % SZ(A)] - A[i]);
                                                             }
for(int i = 0; i < SZ(B); i++)
s2.pb(B[(i + 1) % SZ(B)] - B[i]);
                                                              6.14 Tangent line of Two Circle
 for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
                                                             vector<Line> go(const Cir &c1, const Cir &c2,
 if (p2 >= SZ(B)
                                                                int sign1) {
    || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
                                                               // sign1 = 1 for outer tang, -1 for inter tang
   C.pb(C.back() + s1[p1++]);
                                                               vector<Line> ret;
 else
                                                               if (norm(c1.o - c2.o) < eps)
   C.pb(C.back() + s2[p2++]);
                                                                return ret;
return hull(C), C;
                                                               11f d = abs(c1.o - c2.o);
                                                               Pointf v = (c2.0 - c1.0) / d;
6.11 Circle Class
                                                               11f c = (c1.r - sign1 * c2.r) / d;
                                                               if (c * c > 1)
struct Circle { Pointf o; llf r; };
                                                                return ret;
                                                               11f h = sqrt(max(0.0, 1.0 - c * c));
vector<llf> intersectAngle(Circle A, Circle B) {
                                                               for (int sign2: {1, -1}) {
Pointf n = c * v + sign2 * h * rot90(v);
Pointf dir = B.o - A.o; llf d2 = norm(dir);
if (norm(A.r - B.r) >= d2)
                                                                Pointf p1 = c1.o + n * c1.r;
 if (A.r < B.r) return {-PI, PI}; // special</pre>
                                                                Pointf p2 = c2.0 + n * (c2.r * sign1);
  else return {};
                                                                if (norm(p2 - p1) < eps)
 if (norm(A.r + B.r) <= d2) return {};</pre>
                                                                p2 = p1 + rot90(c2.o - c1.o);
llf dis = abs(dir), theta = arg(dir);
llf phi = acos((A.r * A.r + d2 - B.r * B.r) /
                                                                ret.push_back({p1, p2});
   (2 * A.r * dis));
                                                               return ret;
11f L = theta - phi, R = theta + phi;
                                                             }
while (L < -PI) L += PI * 2;
while (R > PI) R -= PI * 2;
                                                              6.15 Minimum Covering Circle
return { L, R };
                                                              template<typename P>
                                                              Circle getCircum(const P &a, const P &b, const P &c){
                                                               Real a1 = a.x-b.x, b1 = a.y-b.y;
vector<Pointf> intersectPoint(Circle a, Circle b) {
                                                               Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
11f d = abs(a.o - b.o);
                                                               Real a2 = a.x-c.x, b2 = a.y-c.y;
 if (d >= b.r+a.r || d <= abs(b.r-a.r)) return {};</pre>
                                                               Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
llf dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
Pointf dir = (a.o - b.o); dir /= d;
                                                               Circle cc;
                                                               cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
Pointf pcrs = dir*d1 + b.o;
                                                               cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2)
dt=sqrt(max(0.0L, b.r*b.r-d1*d1)), dir = rot90(dir);
                                                               cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
return {pcrs + dir*dt, pcrs - dir*dt};
                                                               return cc;
6.12 Intersection of line and Circle
                                                              template<typename P>
vector<pdd> line_interCircle(const pdd &p1,
                                                              Circle MinCircleCover(const vector<P>& pts){
 const pdd &p2, const pdd &c, const double r) {
                                                               {\tt random\_shuffle(pts.begin(),\ pts.end());}
 pdd ft = p1 + project(c-p1, p2-p1), vec = p2-p1;
                                                               Circle c = { pts[0], 0 };
llf dis = abs(c - ft);
                                                               for(int i=0;i<(int)pts.size();i++){</pre>
if (abs(dis - r) < eps) return {ft};</pre>
                                                                if (dist(pts[i], c.o) <= c.r) continue;</pre>
if (dis > r) return {};
                                                                c = \{ pts[i], 0 \};
vec = vec * sqrt(r * r - dis * dis) / abs(vec);
                                                                for (int j = 0; j < i; j++) {
   if(dist(pts[j], c.o) <= c.r) continue;</pre>
return {ft + vec, ft - vec};
                                                                 c.o = (pts[i] + pts[j]) / 2;
                                                                 c.r = dist(pts[i], c.o);
     Intersection of Polygon and Circle
                                                                 for (int k = 0; k < j; k++) {
                                                                  if (dist(pts[k], c.o) <= c.r) continue;</pre>
// Divides into multiple triangle, and sum up
                                                                  c = getCircum(pts[i], pts[j], pts[k]);
// test by HDU2892
const double PI=acos(-1);
double _area(pdd pa, pdd pb, double r){
                                                                }
if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
if(abs(pb)<eps) return 0;</pre>
                                                               return c;
```

# 6.16 KDTree (Nearest Point) const int MXN = 100005; struct KDTree {

```
struct Node {
  int x,y,x1,y1,x2,y2;
  int id,f;
Node *L, *R;
 } tree[MXN], *root;
 int n:
 LL dis2(int x1, int y1, int x2, int y2) {
 LL dx = x1-x2, dy = y1-y2;
  return dx*dx+dy*dy;
 static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
 static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 void init(vector<pair<int,int>> ip) {
 n = ip.size();
  for (int i=0; i<n; i++) {
   tree[i].id = i;
   tree[i].x = ip[i].first;
   tree[i].y = ip[i].second;
  root = build_tree(0, n-1, 0);
 Node* build_tree(int L, int R, int d) {
  if (L>R) return nullptr
  int M = (L+R)/2; tree[M].f = d%2;
  nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
  tree[M].x1 = tree[M].x2 = tree[M].x;
  tree[M].y1 = tree[M].y2 = tree[M].y;
  tree[M].L = build_tree(L, M-1, d+1);
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
   tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
  tree[M].R = build_tree(M+1, R, d+1);
  if (tree[M].R) {
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
   tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
   tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
  }
  return tree+M;
 int touch(Node* r, int x, int y, LL d2){
 LL dis = sqrt(d2)+1;
  if (x<r->x1-dis || x>r->x2+dis ||
    y<r->y1-dis || y>r->y2+dis)
   return 0;
  return 1:
 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
  if (!r || !touch(r, x, y, md2)) return;
  LL d2 = dis2(r->x, r->y, x, y);
  if (d2 < md2 || (d2 == md2 && mID < r->id)) {
  mID = r -> id;
   md2 = d2;
  }
  // search order depends on split dim
  if ((r->f == 0 \&\& x < r->x) ||
    (r->f == 1 && y < r->y))
   nearest(r->L, x, y, mID, md2);
   nearest(r->R, x, y, mID, md2);
  } else {
   nearest(r->R, x, y, mID, md2);
   nearest(r->L, x, y, mID, md2);
 int query(int x, int y) {
  int id = 1029384756;
  LL d2 = 102938475612345678LL;
  nearest(root, x, y, id, d2);
  return id;
 }
} tree;
```

# 6.17 Rotating Sweep Line

```
void rotatingSweepLine(pair<int, int> a[], int n) {
vector<pair<int, int>> 1;
```

```
1.reserve(n * (n - 1) / 2);
 for (int i = 0; i < n; ++i)
  for (int j = i + 1; j < n; ++j)
   1.emplace_back(i, j);
 sort(1.begin(), 1.end(), [&a](auto &u, auto &v){
  lld udx = a[u.first].first - a[u.second].first;
lld udy = a[u.first].second - a[u.second].second;
  1ld vdx = a[v.first].first - a[v.second].first;
  11d vdy = a[v.first].second - a[v.second].second;
  if (udx == 0 or vdx == 0) return not udx == 0;
  int s = sgn(udx * vdx);
  return udy * vdx * s < vdy * udx * s;</pre>
 });
 vector<int> idx(n), p(n);
 iota(idx.begin(), idx.end(), 0);
 sort(idx.begin(), idx.end(), [&a](int i, int j){
  return a[i] < a[j]; });
 for (int i = 0; i < n; ++i) p[idx[i]] = i;</pre>
 for (auto [i, j]: 1) {
  // do here
  swap(p[i], p[j]);
  idx[p[i]] = i, idx[p[j]] = j;
}
6.18
       Circle Cover
const int N = 1021;
struct CircleCover {
 int C
 Cir c[N]
 bool g[N][N], overlap[N][N];
 // Area[i] : area covered by at least i circles
 double Area[ N ];
 void init(int _C){ C = _C;}
 struct Teve {
  pdd p; double ang; int add;
Teve() {}
  Teve(pdd _a, double _b, int _c):p(_a), ang(_b), add(
     _c){}
  bool operator<(const Teve &a)const
  {return ang < a.ang;}
 }eve[N * 2];
 // strict: x = 0, otherwise x = -1
 bool disjuct(Cir &a, Cir &b, int x)
 {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
bool contain(Cir &a, Cir &b, int x)
 {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
 bool contain(int i, int j) {
  /* c[j] is non-strictly in c[i]. */
  return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c[i].R) \mid c[i].R - c[i].R
     [j].R) == 0 \&\& i < j)) \&\& contain(c[i], c[j], -1);
 void solve(){
  fill_n(Area, C + 2, 0);
  for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)
    overlap[i][j] = contain(i, j);
  for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)
    g[i][j] = !(overlap[i][j] || overlap[j][i] ||
       disjuct(c[i], c[j], -1));
  for(int i = 0; i < C; ++i){
   int E = 0, cnt = 1;
   for(int j = 0; j < C; ++j)
    if(j != i && overlap[j][i])
     ++cnt:
   for(int j = 0; j < C; ++j)
    if(i != j && g[i][j]) {
  pdd aa, bb;
     CCinter(c[i], c[j], aa, bb);
     llf A = atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);
llf B = atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);
     eve[E++] = Teve(bb,B,1), eve[E++]=Teve(aa,A,-1);
     if(B > A) ++cnt;
   if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
   else{
    sort(eve, eve + E);
    eve[E] = eve[0];
    for(int j = 0; j < E; ++j){
     cnt += eve[j].add;
```

Area[cnt] += cross(eve[j].p, eve[j + 1].p) \* .5;

```
double theta = eve[j + 1].ang - eve[j].ang;
    if (theta < 0) theta += 2. * pi;
    Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
}
}
}
}
}</pre>
```

# 7 Stringology

## 7.1 Hash

```
class Hash {
  private:
    static constexpr int P = 127, Q = 1051762951;
  vector<int> h, p;
  public:
    void init(const string &s){
      h.assign(s.size()+1, 0); p.resize(s.size()+1);
      for (size_t i = 0; i < s.size(); ++i)
          h[i + 1] = add(mul(h[i], P), s[i]);
      generate(p.begin(), p.end(), [x=1,y=1,this]()
          mutable{y=x;x=mul(x,P);return y;});
    }
    int query(int l, int r){ // 1-base (1, r]
      return sub(h[r], mul(h[1], p[r-1]));}
};</pre>
```

```
7.2 Suffix Array
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
int x[maxn], p[maxn], q[maxn * 2];
// sa[i]: sa[i]-th suffix is the \
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix \
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *sa, int *c, int n, int z) {
memset(sa, 0, sizeof(int) * n);
 memcpy(x, c, sizeof(int) * z);
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
 if (sa[i] && !t[sa[i] - 1])
   sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
 memcpy(x, c, sizeof(int) * z);
 for (int i = n - 1; i >= 0; --i)
  if (sa[i] && t[sa[i] - 1])
   sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q,
bool *t, int *c, int n, int z) {
bool uniq = t[n - 1] = true;
 int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
 memset(c, 0, sizeof(int) * z);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
 for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
 if (uniq) {
 for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
  return:
 for (int i = n - 2; i \ge 0; --i)
 t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 pre(sa, c, n, z);
 for (int i = 1; i <= n - 1; ++i)
if (t[i] && !t[i - 1])
   sa[--x[s[i]]] = p[q[i] = nn++] = i;
 induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) {
  if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
  bool neq = last < 0 || \</pre>
   memcmp(s + sa[i], s + last,
(p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
  ns[q[last = sa[i]]] = nmxz += neq;
 }}
 sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
 pre(sa, c, n, z);
for (int i = nn - 1; i >= 0; --i)
  sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
```

```
induce(sa, c, s, t, n, z);
void build(const string &s) {
 for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
 _s[(int)s.size()] = 0; // s shouldn't contain 0
 sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];
for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
 int ind = 0; hi[0] = 0;
 for (int i = 0; i < (int)s.size(); ++i) {
  if (!rev[i]) {</pre>
   ind = 0;
   continue:
  while (i + ind < (int)s.size() && \</pre>
   s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  hi[rev[i]] = ind ? ind-- : 0;
7.3 Suffix Automaton
struct SuffixAutomaton {
 struct node {
  int ch[K], len, fail, cnt, indeg;
  node(int L = 0) : ch{}, len(L), fail(0), cnt(0),
     indea(0) {}
   st[N];
```

```
int root, last, tot;
void extend(int c) {
 int cur = ++tot;
 st[cur] = node(st[last].len + 1);
 while (last && !st[last].ch[c]) {
   st[last].ch[c] = cur;
   last = st[last].fail;
 if (!last) {
   st[cur].fail = root;
 } else {
   int q = st[last].ch[c];
   if (st[q].len == st[last].len + 1) {
     st[cur].fail = q;
   } else {
     int clone = ++tot;
     st[clone] = st[q];
     st[clone].len = st[last].len + 1;
     st[st[cur].fail = st[q].fail = clone].cnt = 0;
     while (last && st[last].ch[c] == q) {
       st[last].ch[c] = clone;
       last = st[last].fail;
     }
  }
st[last = cur].cnt += 1;
void init(const char* s) {
 root = last = tot = 1;
 st[root] = node(0);
 for (char c; c = *s; ++s) extend(c - 'a');
int q[N];
void dp() +
 for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
 int head = 0, tail = 0;
 for (int i = 1; i <= tot; i++)
   if (st[i].indeg == 0) q[tail++] = i;
 while (head != tail) {
   int now = q[head++];
   if (int f = st[now].fail) {
     st[f].cnt += st[now].cnt;
     if (--st[f].indeg == 0) q[tail++] = f;
int run(const char* s) {
 int now = root;
for (char c; c = *s; ++s) {
   if (!st[now].ch[c -= 'a']) return 0;
   now = st[now].ch[c];
 return st[now].cnt;
```

```
| } SAM;
 7.4 KMP
 vector<int> kmp(const string &s) {
 vector<int> f(s.size(), 0);
 /* f[i] = length of the longest prefix
    (excluding s[0:i]) such that it coincides
    with the suffix of s[0:i] of the same length */
  /* i + 1 - f[i] is the length of the
    smallest recurring period of s[0:i] */
  int k = 0:
 for (int i = 1; i < (int)s.size(); ++i) {</pre>
   while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
   if (s[i] == s[k]) ++k;
   f[i] = k;
 return f;
 vector<int> search(const string &s, const string &t) {
 // return 0-indexed occurrence of t in s
  vector<int> f = kmp(t), r;
 for (int i = 0, k = 0; i < (int)s.size(); ++i) {
   while(k > 0 \& (k==(int)t.size() \mid \mid s[i]!=t[k]))
    k = f[k - 1];
   if (s[i] == t[k]) ++k;
   if (k == (int)t.size()) r.push_back(i-t.size()+1);
 return res;
 7.5 Z value
 vector<int> Zalgo(const string &s) {
 vector<int> z(s.size(), s.size());
 for (int i = 1, 1 = 0, r = 0; i < z[0]; ++i) {
  int j = clamp(r - i, 0, z[i - 1]);</pre>
   for (; i + j < z[0] \text{ and } s[i + j] == s[j]; ++j);
   if (i + (z[i] = j) > r) r = i + z[1 = i];
 return z;
}
 7.6 Manacher
 int z[maxn];
int manacher(const string& s) {
  string t = ".";
  for(char c: s) t += c, t += '.';
 int 1 = 0, r = 0, ans = 0;
 for (int i = 1; i < t.length(); ++i) {
  z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
  while (i - z[i] >= 0 && i + z[i] < t.length()) {</pre>
    if(t[i - z[i]] == t[i + z[i]]) ++z[i];
    else break;
   if (i + z[i] > r) r = i + z[i], l = i;
 for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
 return ans;
 7.7 Lexico Smallest Rotation
 string mcp(string s) {
 int n = s.length();
  s += s; int i = 0, j = 1;
 while (i < n && j < n) {</pre>
   int k = 0;
   while (k < n \&\& s[i + k] == s[j + k]) k++;
   ((s[i+k] \leftarrow s[j+k]) ? j : i) += k + 1;
   j += (i == j);
 return s.substr(i < n ? i : j, n);</pre>
 7.8 BWT
 struct BurrowsWheeler{
```

#define SIGMA 26

#define BASE 'a

vector<int> v[ SIGMA ];

// make ori -> ori + ori
// then build suffix array

void BWT(char\* ori, char\* res){

```
void iBWT(char* ori, char* res){
  for( int i = 0 ; i < SIGMA ; i ++ )
    v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )
    v[ ori[i] - BASE ].push_back( i );
  vector<int> a;
  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )
    for( auto j : v[ i ] ){
      a.push_back( j );
      ori[ ptr ++ ] = BASE + i;
    }
  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
    res[ i ] = ori[ a[ ptr ] ];
    ptr = a[ ptr ];
  }
  res[ len ] = 0;
}
bwt;</pre>
```

#### 7.9 Palindromic Tree

```
struct palindromic_tree{
 struct node{
  int next[26],f,len;
  int cnt,num,st,ed;
  node(int l=0):f(0),len(1),cnt(0),num(0) {
   memset(next, 0, sizeof(next)); }
 vector<node> st;
 vector<char> s;
 int last,n;
 void init(){
  st.clear();s.clear();last=1; n=0;
  st.push_back(0);st.push_back(-1);
  st[0].f=1;s.push_back(-1); }
 int getFail(int x){
  while(s[n-st[x].len-1]!=s[n])x=st[x].f;
  return x;}
 void add(int c){
  s.push_back(c-='a'); ++n;
  int cur=getFail(last);
  if(!st[cur].next[c]){
   int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
   st[now].num=st[st[now].f].num+1;
  last=st[cur].next[c];
  ++st[last].cnt;}
 void dpcnt() {
  for (int i=st.size()-1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
 int size(){ return st.size()-2;}
} pt;
int main() {
 string s; cin >> s; pt.init();
 for (int i=0; i<SZ(s); i++) {
  int prvsz = pt.size(); pt.add(s[i]);</pre>
  if (prvsz != pt.size()) {
   int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [1,r]: s.substr(1, r-l+1)
 }
 return 0;
```

### 8 Misc

# 8.1 Theorems

# 8.1.1 Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii}=d(i)$ ,  $L_{ij}=-c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|{\rm det}(\tilde{L}_{11})|.$
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .

#### 8.1.2 Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij}=x_{ij}$  ( $x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij}=-d_{ji}.$   $\frac{rank(D)}{2}$  is the maximum matching on C

### 8.1.3 Cayley's Formula

- Given a degree sequence  $d_1,d_2,\dots,d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .

#### 8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \ge d_2 \ge \ldots \ge d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1 + d_2 + \ldots + d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all  $1 \le k \le n$ .

#### 8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

#### 8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let  $N_G(W)$  denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff  $\forall W\subseteq X, |W|\leq |N_G(W)|$ 

#### 8.1.7 Euler's planar graph formula

$$V - E + F = C + 1, E \le 3V - 6$$
(?)

# 8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have  $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$ 

#### 8.1.9 Lucas's theorem

 ${m\choose n}\equiv\prod_{i=0}^k{m_i\choose n_i}\pmod{p}, \text{ where } m=m_kp^k+m_{k-1}p^{k-1}+\cdots+m_1p+m_0,$  and  $n=n_kp^k+n_{k-1}p^{k-1}+\cdots+n_1p+n_0.$ 

#### 8.1.10 Matroid Intersection

Given matroids  $M_1=(G,I_1),M_2=(G,I_2)$ , find maximum  $S\in I_1\cap I_2$ . For each iteration, build the directed graph and find a shortest path from s to t.

- $s \rightarrow x : S \sqcup \{x\} \in I_1$
- $x \to t : S \sqcup \{x\} \in I_2$
- $y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1$  (y is in the unique circuit of  $S \sqcup \{x\}$ )
- $x \to y: S \setminus \{y\} \sqcup \{x\} \in I_2$  (y is in the unique circuit of  $S \sqcup \{x\}$ )

Alternate the path, and |S| will increase by 1. Let  $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)), N=|G|$ . In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to  $x\in S$  and  $x\notin S$ , resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

# 8.2 DP-opt Condition

# 8.2.1 totally monotone (concave/convex)

 $\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}$ 

#### 8.2.2 monge condition (concave/convex)

 $\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}$ 

## 8.3 Convex 1D/1D DP

```
struct segment {
  int i, 1, r;
  segment() {}
  segment(int a, int b, int c): i(a), l(b), r(c) {}
};
  inline lld f(int l, int r){return dp[l] + w(l+1, r);}
void solve() {
  dp[0] = 0;
  deque<segment> dq; dq.push_back(segment(0, 1, n));
  for (int i = 1; i <= n; ++i) {
    dp[i] = f(dq.front().i, i);
}</pre>
```

```
while(dq.size()&&dq.front().r<i+1) dq.pop_front();
dq.front().l = i + 1;
segment seg = segment(i, i + 1, n);
while (dq.size() &&
    f(i, dq.back().l)<f(dq.back().i, dq.back().l))
        dq.pop_back();
if (dq.size()) {
    int d = 1 << 20, c = dq.back().l;
    while (d >>= 1) if (c + d <= dq.back().r)
        if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
        dq.back().r = c; seg.l = c + 1;
    }
    if (seg.l <= n) dq.push_back(seg);
}
</pre>
```

# 8.4 ConvexHull Optimization

```
struct L {
mutable int64_t a, b, p;
 bool operator<(const L &r) const { return a < r.a; }</pre>
bool operator<(int64_t x) const { return p < x; }</pre>
struct DynamicHull : multiset<L, less<>> {
 static const int64_t kInf = 1e18;
 bool Isect(iterator x, iterator y)
  auto Div = [](int64_t a, int64_t b) {
    return a / b - ((a ^ b) < 0 && a % b); }
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(int64_t a, int64_t b) {
 auto z = insert({a, b, 0}), y = z++, x = y;
while (Isect(y, z)) z = erase(z);
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
  while ((y = x) != begin() \&\& (--x)->p >= y->p)
   Isect(x, erase(y));
 int64_t Query(int64_t x) {
 auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
```

# 8.5 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

## 8.6 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
 dfn[u]=low[u]=++dfs_idx;
 for(int i=0;i<(int)init_g[u].size();i++){</pre>
  int v=init_g[u][i];
  if(v==par[u]) continue;
  if(!dfn[v]){
   par[v]=u;
   tarjan(v);
   low[u]=min(low[u],low[v]);
   if(dfn[u]<low[v]){</pre>
    g[u].push_back(v)
    g[v].push_back(u);
  }else{
   low[u]=min(low[u],dfn[v]);
   if(dfn[v]<dfn[u]){</pre>
    int temp_v=u;
    bcc_id++;
```

```
while(temp_v!=v){
                                                                   for(int i=0;i<=mx-obj[s].FF;i++)</pre>
                                                                    dp[s][i] = dp[u][i];
     g[bcc_id+n].push_back(temp_v);
                                                                   dfs(s, mx - obj[s].first);
     g[temp_v].push_back(bcc_id+n);
     temp_v=par[temp_v];
                                                                   for(int i=obj[s].FF;i<=mx;i++)</pre>
                                                                    dp[u][i] = max(dp[u][i],
    g[bcc_id+n].push_back(v);
                                                                      dp[s][i - obj[s].FF] + obj[s].SS);
    g[v].push_back(bcc_id+n);
                                                                 }
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
                                                                 8.8 N Queens Problem
                                                                 vector< int > solve( int n ) {
                                                                  // no solution when n=2, 3
int dp[maxn][2],min_dp[2][2],tmp[2][2],tp[2];
                                                                  vector< int > ret;
void dfs(int u,int fa){
                                                                  if ( n % 6 == 2 ) {
  for ( int i = 2 ; i <= n ; i += 2 )
    ret.push_back( i );</pre>
 if(u<=n){
  for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                   ret.push_back( 3 ); ret.push_back( 1 );
for ( int i = 7 ; i <= n ; i += 2 )
ret.push_back( i );</pre>
   int v=g[u][i];
   if(v==fa) continue;
   dfs(v,u);
   memset(tp,0x8f,sizeof tp);
                                                                   ret.push_back( 5 );
   if(v<=n){
                                                                  } else if ( n % 6 == 3 ) {
for ( int i = 4 ; i <= n ; i += 2 )</pre>
    tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
    tp[1]=max(
                                                                    ret.push_back( i );
     dp[u][0]+dp[v][0]+1
                                                                   ret.push_back( 2 );
     dp[u][1]+max(dp[v][0],dp[v][1])
                                                                   for ( int i = 5 ; i <= n ; i += 2 )
                                                                    ret.push_back( i );
   }else{
                                                                   ret.push_back( 1 ); ret.push_back( 3 );
    tp[0]=dp[u][0]+dp[v][0];
                                                                  } else {
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
                                                                   for ( int i = 2 ; i <= n ; i += 2 )
                                                                    ret.push_back( i );
   dp[u][0]=tp[0],dp[u][1]=tp[1];
                                                                   for ( int i = 1 ; i <= n ; i += 2 )
                                                                    ret.push_back( i );
 }else{
  for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                  return ret;
   int v=g[u][i];
   if(v==fa) continue;
   dfs(v,u);
                                                                 8.9 Aliens Optimization
                                                                 long long Alien() {
  min_dp[0][0]=0;
                                                                  long long c = kInf;
  min_dp[1][1]=1;
                                                                  for (int d = 60; d >= 0; --d) {
  min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
                                                                   // cost can be negative, depending on the problem. if (c - (1LL << d) < 0) continue;
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
                                                                   long long ck = c - (1LL \ll d);
   if(v==fa) continue;
                                                                   pair<long long, int> r = check(ck);
   memset(tmp,0x8f,sizeof tmp);
                                                                   if (r.second == k) return r.first - ck * k;
   tmp[0][0]=max(
                                                                   if (r.second < k) c = ck;</pre>
    \min_{dp[0][0]+\max(dp[v][0],dp[v][1])}
    min_dp[0][1]+dp[v][0]
                                                                  pair<long long, int> r = check(c);
return r.first - c * k;
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
   tmp[1][0]=max(
    \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
                                                                        Hilbert Curve
                                                                 8.10
    min_dp[1][1]+dp[v][0]
                                                                 long long hilbert(int n, int x, int y) {
                                                                  long long res = 0;
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
                                                                  for (int s = n / 2; s; s >>= 1) {
   memcpy(min_dp,tmp,sizeof tmp);
                                                                   int rx = (x & s) > 0, ry = (y & s) > 0;
res += s * 111 * s * ((3 * rx) ^ ry);
  dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
                                                                   if (ry == 0) {
  dp[u][0]=min_dp[0][0];
                                                                    if (rx == 1) x = s - 1 - x, y = s - 1 - y;
                                                                     swap(x, y);
int main(){
 int m,a,b;
                                                                  return res;
 scanf("%d%d",&n,&m);
 for(int i=0;i<m;i++){
  scanf("%d%d",&a,&b);</pre>
                                                                 8.11 Binary Search On Fraction
  init_g[a].push_back(b);
  init_g[b].push_back(a);
                                                                 struct 0 {
                                                                  11 p, q;
 par[1]=-1;
                                                                  Q go(Q b, 11 d) \{ return \{ p + b.p*d, q + b.q*d \}; \}
 tarjan(1);
 dfs(1,-1);
                                                                 bool pred(Q);
 printf("%d\n", max(dp[1][0], dp[1][1]));
                                                                 // returns smallest p/q in [lo, hi] such that
 return 0;
                                                                 // pred(p/q) is true, and 0 <= p,q <= N
                                                                 Q frac_bs(11 N) {
                                                                  Q lo{0, 1}, hi{1, 0};
     Tree Knapsack
                                                                  if (pred(lo)) return lo;
int dp[N][K]; PII obj[N];
                                                                  assert(pred(hi));
vector<int> G[N];
                                                                  bool dir = 1, L = 1, H = 1;
                                                                  for (; L || H; dir = !dir) {
void dfs(int u, int mx){
for(int s: G[u]) {
                                                                   11 len = 0, step = 1;
  if(mx < obj[s].first) continue;</pre>
                                                                   for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
```

```
if (Q mid = hi.go(lo, len + step);
    mid.p > N || mid.q > N || dir ^ pred(mid))
    t++;
    else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
}
return dir ? hi : lo;
}
```