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8.2 MaximumEmptyRect

8.3 DP-opt Condition
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8.4 Convex ID/ID DP

8.5 ConvexHull Optimization

8.6 Josephus Problem

8.7 Cactus Matching

8.8 DLX
  Basic
    vimrc
se is nu rnu bs=2 ru mouse=a encoding=utf-8
se cin et sw=4 sts=4 t_Co=256 tgc sc hls ls=2
```

```
syn on
colorscheme desert
filetype indent on
inoremap {<CR>} {<CR>} < ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 -
    DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
    Wconversion -fsanitize=address -fsanitize=undefined
     -g && echo success<CR>
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 -
    DKISEKI && echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
```

## 1.2 Increase Stack

```
const int size = 256 << 20;</pre>
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));
```

### 1.3 Pragma Optimization

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,sse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
```

### 1.4 IO Optimization

```
static inline int gc() {
  static char buf[ 1 << 20 ], *p = buf, *end = buf;</pre>
 if ( p == end ) {
  end = buf + fread( buf, 1, 1 << 20, stdin );
  if ( end == buf ) return EOF;
  p = buf;
 return *p++;
template < typename T >
static inline bool gn( T &_ ) {
  register int c = gc(); register T __
 while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
 if(c == '-') { __ = -1; c = gc(); }
if(c == EOF) return false;
 while('0' <= c\&c <= '9') _ = _ * 10 + c - '0', c = gc();
 _ *= __;
 return true;
```

```
template < typename T, typename ...Args >
static inline bool gn( T &x, Args &...args )
{ return gn(x) && gn(args...); }
```

## 2 Data Structure

## 2.1 Bigint

```
class BigInt{
private:
using lld = int_fast64_t;
#define PRINTF_ARG PRIdFAST64
#define LOG_BASE_STR "9"
static constexpr lld BASE = 1000000000;
static constexpr int LOG_BASE = 9;
vector<lld> dig; bool neg;
inline int len() const { return (int) dig.size(); }
inline int cmp_minus(const BigInt& a) const {
 if(len() == 0 && a.len() == 0) return 0;
  if(neg ^ a.neg)return a.neg ^ 1;
 if(len()!=a.len())
   return neg?a.len()-len():len()-a.len();
 for(int i=len()-1;i>=0;i--) if(dig[i]!=a.dig[i])
  return neg?a.dig[i]-dig[i]:dig[i]-a.dig[i];
 return 0;
inline void trim(){
 while(!dig.empty()&&!dig.back())dig.pop_back();
 if(dig.empty()) neg = false;
}
public:
BigInt(): dig(vector<lld>()), neg(false){}
BigInt(lld a): dig(vector<lld>()){
 neg = a<0; dig.push_back(abs(a));</pre>
 trim();
BigInt(const string& a): dig(vector<lld>()){
 assert(!a.empty()); neg = (a[0]=='-');
  for(int i=((int)a.size())-1;i>=neg;i-=LOG_BASE){
  11d cur = 0;
  for(int j=min(LOG_BASE-1,i-neg);j>=0;j--)
    cur = cur*10+a[i-j]-'0';
  dig.push_back(cur);
 } trim();
inline bool operator<(const BigInt& a)const</pre>
  {return cmp_minus(a)<0;}
 inline bool operator <= (const BigInt& a) const
  {return cmp_minus(a)<=0;}
inline bool operator == (const BigInt& a)const
 {return cmp_minus(a)==0;}
inline bool operator!=(const BigInt& a)const
  {return cmp_minus(a)!=0;}
 inline bool operator>(const BigInt& a)const
  {return cmp_minus(a)>0;}
 inline bool operator>=(const BigInt& a)const
  {return cmp_minus(a)>=0;}
BigInt operator-() const {
 BigInt ret = *this;
 ret.neg ^= 1; return ret;
BigInt operator+(const BigInt& a) const {
 if(neg) return -(-(*this)+(-a));
  if(a.neg) return (*this)-(-a);
  int n = max(a.len(), len())
 BigInt ret; ret.dig.resize(n);
 11d pro = 0;
  for(int i=0;i<n;i++) {</pre>
   ret.dig[i] = pro;
  if(i < a.len()) ret.dig[i] += a.dig[i];</pre>
  if(i < len()) ret.dig[i] += dig[i];</pre>
  pro = 0;
   if(ret.dig[i] >= BASE) pro = ret.dig[i]/BASE;
   ret.dig[i] -= BASE*pro;
 if(pro != 0) ret.dig.push_back(pro);
 return ret;
BigInt operator-(const BigInt& a) const {
 if(neg) return -(-(*this) - (-a));
  if(a.neg) return (*this) + (-a);
  int diff = cmp_minus(a);
 if(diff < 0) return -(a - (*this));</pre>
```

```
if(diff == 0) return 0;
  BigInt ret; ret.dig.resize(len(), 0);
  for(int i=0;i<len();i++) {</pre>
   ret.dig[i] += dig[i];
   if(i < a.len()) ret.dig[i] -= a.dig[i];</pre>
   if(ret.dig[i] < 0){</pre>
    ret.dig[i] += BASE;
    ret.dig[i+1]--;
  ret.trim(); return ret;
 BigInt operator*(const BigInt& a) const {
  if(!len()||!a.len()) return 0;
  BigInt ret; ret.dig.resize(len()+a.len()+1);
  ret.neg = neg ^ a.neg;
  for(int i=0;i<len();i++)</pre>
   for(int j=0;j<a.len();j++){</pre>
    ret.dig[i+j] += dig[i] * a.dig[j];
    if(ret.dig[i+j] >= BASE) {
     lld x = ret.dig[i+j] / BASE;
     ret.dig[i+j+1] += x;
     ret.dig[i+j] -= x * BASE;
  ret.trim(); return ret;
 BigInt operator/(const BigInt& a) const {
  assert(a.len());
  if(len() < a.len()) return 0;</pre>
  BigInt ret; ret.dig.resize(len()-a.len()+1);
  ret.neg = a.neg;
  for(int i=len()-a.len();i>=0;i--){
   11d 1 = 0, r = BASE;
   while(r-l > 1){
    lld mid = (1+r)>>1;
ret.dig[i] = mid;
    if(ret*a<=(neg?-(*this):(*this))) 1 = mid;</pre>
    else r = mid;
   ret.dig[i] = 1;
  ret.neg ^= neg; ret.trim();
  return ret;
 BigInt operator%(const BigInt& a) const {
  return (*this) - (*this) / a * a;
 friend BigInt abs(BigInt a) { a.neg = 0; return a; }
 friend void swap(BigInt& a, BigInt& b){
  swap(a.dig, b.dig); swap(a.neg, b.neg);
 friend istream& operator>>(istream& ss, BigInt& a){
  string s; ss >> s; a = s; return ss;
 friend ostream&operator<<(ostream&o, const BigInt&a){</pre>
  if(a.len() == 0) return o << '0';
if(a.neg) o << '-';</pre>
  if(a.neg) o <<</pre>
  o << a.dig.back();
  for(int i=a.len()-2;i>=0;i--)
   o<<setw(LOG_BASE)<<setfill('0')<<a.dig[i];
  return o;
 inline void print() const {
  if(len() == 0){putchar('0');return;}
  if(neg) putchar('-');
printf("%" PRINTF_ARG, dig.back());
  for(int i=len()-2;i>=0;i--
   printf("%0" LOG_BASE_STR PRINTF_ARG, dig[i]);
 #undef PRINTF_ARG
 #undef LOG_BASE_STR
2.2 Dark Magic
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using __gnu_pbds::pairing_heap_tag;
using __gnu_pbds::binary_heap_tag;
using __gnu_pbds::binomial_heap_tag;
using __gnu_pbds::rc_binomial_heap_tag;
```

using \_\_gnu\_pbds::thin\_heap\_tag;

```
template<typename T>
                                                               return par==nullptr ||\
using pbds_heap=__gnu_pbds::prioity_queue<T, less<T>, \
                                                                (par->ch[0]!=this && par->ch[1]!=this);
                    pairing_heap_tag>;
// __gnu_pbds::priority_queue<T,less<T>>::
                                                              bool is_rch(){return !is_root() && par->ch[1]==this;}
                                                             } *node[maxn], *stk[maxn];
    point_iterator
                                                             int top;
// x = pq.push(10); pq.modify(x, 87); a.join(b);
using __gnu_pbds::rb_tree_tag;
                                                             void to_child(Node* p,Node* c,bool dir){
using __gnu_pbds::ov_tree_tag;
                                                              p->ch[dir]=c;
using __gnu_pbds::splay_tree_tag;
                                                              p->up();
template<typename T>
                                                             inline void rotate(Node* node){
using ordered_set = __gnu_pbds::tree<T,\</pre>
__gnu_pbds::null_type,less<T>,rb_tree_tag,\
                                                              Node* par=node->par;
 _gnu_pbds::tree_order_statistics_node_update>;
                                                              Node* par_par=par->par
// find_by_order, order_of_key
                                                              bool dir=node->is_rch();
template<typename A, typename B>
                                                              bool par_dir=par->is_rch();
using hTable1=__gnu_pbds::cc_hash_table<A,B>;
                                                              to_child(par, node->ch[!dir], dir);
template<typename A, typename B>
                                                              to_child(node,par,!dir);
using hTable2=__gnu_pbds::gp_hash_table<A,B>;
                                                              if(par_par!=nullptr && par_par->ch[par_dir]==par)
                                                               to_child(par_par,node,par_dir);
2.3 Disjoint Set
                                                              else node->par=par_par;
class DJS {
private:
                                                             inline void splay(Node* node){
vector< int > fa, sz, sv;
vector< pair< int*, int > > opt;
                                                              Node* tmp=node;
                                                              stk[top++]=node:
void assign( int *k, int v ) {
                                                              while(!tmp->is_root()){
 opt.emplace_back( k, *k );
                                                               tmp=tmp->par;
                                                               stk[top++]=tmp;
  *k = v;
public:
                                                              while(top) stk[--top]->down();
void init( int n ) {
                                                              for(Node *fa=node->par;
 fa.resize( n ); iota( fa.begin(), fa.end(), 0 );
sz.resize( n ); fill( sz.begin(), sz.end(), 1 );
                                                               !node->is_root();
                                                               rotate(node),fa=node->par)
  opt.clear();
                                                               if(!fa->is_root())
                                                                rotate(fa->is_rch()==node->is_rch()?fa:node);
int query(int x) {return fa[x] == x?x:query(fa[x]);}
void merge( int a, int b ) {
                                                             inline void access(Node* node){
                                                              Node* last=nullptr;
 int af = query( a ), bf = query( b );
                                                              while(node!=nullptr){
  if( af == bf ) return;
 if( sz[ af ] < sz[ bf ] ) swap( af, bf );
assign( &fa[ bf ], fa[ af ] );</pre>
                                                               splay(node);
                                                               to_child(node, last, true);
 assign( &sz[ af ], sz[ af ] + sz[ bf ] );
                                                               last=node;
                                                               node=node->par;
 void save() { sv.push_back( (int) opt.size() ); }
void undo() {
 int ls = sv.back(); sv.pop_back();
                                                             inline void change_root(Node* node){
 while ( ( int ) opt.size() > ls )
                                                              access(node);splay(node);node->set_rev();
  pair< int*, int > cur = opt.back();
                                                             inline void link(Node* x,Node* y){
   *cur.first = cur.second;
                                                              change_root(x);splay(x);x->par=y;
   opt.pop_back();
                                                             inline void split(Node* x,Node* y){
                                                              change_root(x);access(y);splay(x);
                                                              to_child(x,nullptr,true);y->par=nullptr;
2.4 Link-Cut Tree
                                                             inline void change_val(Node* node,int v){
struct Node{
Node *par, *ch[2];
                                                              access(node);splay(node);node->v=v;node->up();
 int xor_sum, v;
bool is_rev;
                                                             inline int query(Node* x, Node* y){
Node(int _v){
                                                              change_root(x);access(y);splay(y);
 v=xor_sum=_v;is_rev=false;
                                                              return y->xor_sum;
 par=ch[0]=ch[1]=nullptr;
                                                             inline Node* find_root(Node* node){
                                                              access(node);splay(node);
inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
                                                              Node* last=nullptr
inline void down(){
  if(is_rev){
                                                              while(node!=nullptr){
   if(ch[0]!=nullptr) ch[0]->set_rev();
                                                               node->down();last=node;node=node->ch[0];
   if(ch[1]!=nullptr) ch[1]->set_rev();
   is_rev=false;
                                                              return last;
  }
                                                             set<pii> dic;
 inline void up(){
                                                             inline void add_edge(int u,int v){
                                                              if(u>v) swap(u,v)
 xor_sum=v;
  if(ch[0]!=nullptr){
                                                              if(find_root(node[u])==find_root(node[v])) return;
  xor_sum^=ch[0]->xor_sum;
                                                              dic.insert(pii(u,v))
                                                              link(node[u],node[v]);
   ch[0]->par=this;
  if(ch[1]!=nullptr){
                                                             inline void del_edge(int u,int v){
  xor_sum^=ch[1]->xor_sum;
                                                              if(u>v) swap(u,v);
   ch[1]->par=this;
                                                              if(dic.find(pii(u,v))==dic.end()) return;
                                                              dic.erase(pii(u,v))
                                                              split(node[u],node[v]);
inline bool is_root(){
```

## 2.5 LiChao Segment Tree

```
struct Line{
int m, k, id;
Line() : id( -1 ) {}
Line( int a, int b, int c )
: m( a ), k( b ), id( c ) {}
int at( int x ) { return m * x + k; }
class LiChao {
private:
 int n; vector< Line > nodes;
 inline int lc( int x ) { return 2 * x + 1; }
inline int rc( int x ) { return 2 * x + 2; }
 void insert( int 1, int r, int id, Line ln ) {
   int m = ( 1 + r ) >> 1;
   if ( nodes[ id ].id == -1 ) {
   nodes[ id ] = ln;
   bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
   if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
    atLeft ^= 1; swap( nodes[ id ], ln );
   if ( r - 1 == 1 ) return;
   if ( atLeft ) insert( 1, m, lc( id ), ln );
   else insert( m, r, rc( id ), ln );
  int query( int 1, int r, int id, int x ) {
   int ret = 0;
   if ( nodes[ id ].id != -1 )
    ret = nodes[ id ].at( x );
   int m = (1 + r) >> 1;
   if ( r - l == 1 ) return ret;
   else if ( x < m )</pre>
    return max( ret, query( 1, m, lc( id ), x ) );
    return max( ret, query( m, r, rc( id ), x ) );
public:
 void build( int n_ ) {
  n = n_; nodes.clear();
   nodes.resize( n << 2, Line() );</pre>
 void insert( Line ln ) { insert( 0, n, 0, ln ); }
 int query( int x ) { return query( 0, n, 0, x ); }
} lichao;
```

## 2.6 Treap

```
namespace Treap{
#define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
struct node{
 int size;
 uint32_t pri;
 node *lc, *rc;
node() : size(0), pri(rand()), lc( 0 ), rc( 0 ) {}
 void pull() {
  size = 1;
  if ( lc ) size += lc->size;
   if ( rc ) size += rc->size;
 }
node* merge( node* L, node* R ) {
 if ( not L or not R ) return L ? L : R;
 if ( L->pri > R->pri ) {
  L->rc = merge( L->rc, R ); L->pull();
  return L;
 } else {
  R->lc = merge( L, R->lc ); R->pull();
   return R;
 }
}
void split_by_size( node*rt,int k,node*&L,node*&R ) {
 if ( not rt ) L = R = nullptr;
  else if( sz( rt->lc ) + 1 <= k ) {
  split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
  L->pull();
 } else {
   split_by_size( rt->lc, k, L, R->lc );
  R->pull();
```

```
}
 #undef sz
     Sparse Table
template < typename T, typename Cmp_ = less< T > >
class SparseTable {
private:
 vector< vector< T > > tbl;
 vector< int > lg;
 T cv(Ta, Tb) {
  return Cmp_()( a, b ) ? a : b;
public:
 void init( T arr[], int n ) {
  // 0-base
  lg.resize(n+1);
  lg[0] = -1;
  for( int i=1 ; i<=n ; ++i ) lg[i] = lg[i>>1] + 1;
  tbl.resize(lg[n] + 1);
  tbl[ 0 ].resize( n );
  copy( arr, arr + n, tbl[ 0 ].begin() );
for ( int i = 1 ; i <= lg[ n ] ; ++ i ) {</pre>
   int len = 1 << ( i - 1 ), sz = 1 << i;
   tbl[ i ].resize( n - sz + 1 );
for ( int j = 0 ; j <= n - sz
                      j <= n - sz ;
    tbl[i][j] = cv(tbl[i-1][j], tbl[i-1][j+len]);
 T query( int 1, int r ) {
  // 0-base [1, r)
  int wh = \lg[r-1],
                         len = 1 << wh;
  return cv( tbl[ wh ][ 1 ], tbl[ wh ][ r - len ] );
};
2.8 Linear Basis
struct LinearBasis {
private:
 int n, sz;
 vector< llu > B;
 inline llu two( int x ){ return ( ( llu ) 1 ) << x; }</pre>
public:
 void init( int n_ ) {
  n = n_{;} B.clear(); B.resize(n); sz = 0;
 void insert( llu x ) {
  // add x into B
  for ( int i = n-1; i >= 0; --i ) if( two(i) & x ){
   if (B[i]) x ^= B[i];
    B[i] = x; sz++;
    for ( int j = i - 1 ; j >= 0 ; -- j )
     if( B[ j ] && ( two( j ) & B[ i ] ) )
B[ i ] ^= B[ j ];
    for (int j = i + 1 ; j < n ; ++ j )
     if ( two( i ) & B['j'] )
B[ j ] ^= B[ i ];
    break;
   }
  }
 inline int size() { return sz; }
 bool check( llu x )
  // is x in span(B) ?
  for ( int i = n-1 ; i >= 0 ; --i ) if( two(i) & x )
   if( B[ i ] ) x ^= B[ i ];
   else return false;
  return true;
 llu kth_small(llu k) {
  /** 1-base would always > 0 **/
  /** should check it **/
  /* if we choose at least one element
    but size(B)(vectors in B)==N(original elements)
    then we can't get 0 */
  11u ret = 0;
  for ( int i = 0 ; i < n ; ++ i ) if( B[ i ] ) {
   if( k & 1 ) ret ^= B[ i ];
```

k >>= 1;

ap[u] = true;

```
while (true) {
  return ret;
                                                                     int eid = st.back(); st.pop_back();
} base;
                                                                     bcc[eid] = ecnt;
                                                                     if (eid == t) break;
     Graph
3
                                                                    ecnt++;
    Euler Circuit
bool vis[ N ]; size_t la[ K ];
                                                                  if (ch == 1 and u == f) ap[u] = false;
void dfs( int u, vector< int >& vec ) {
while ( la[ u ] < G[ u ].size() ) {</pre>
                                                                public:
  if( vis[ G[ u ][ la[ u ] ].second ] ) {
                                                                 void init(int n_) {
   ++ la[ u ];
                                                                 G.clear(); G.resize(n = n_);
   continue:
                                                                  ecnt = 0; ap.assign(n, false);
                                                                  low.assign(n, 0); dfn.assign(n, 0);
 int v = G[ u ][ la[ u ] ].first;
vis[ G[ u ][ la[ u ] ].second ] = true;
                                                                 void add_edge(int u, int v) {
  ++ la[ u ]; dfs( v, vec );
                                                                 G[u].emplace_back(v, ecnt);
  vec.push_back( v );
                                                                 G[v].emplace_back(u, ecnt++);
                                                                 void solve() {
                                                                  ins.assign(ecnt, false);
3.2 BCC Edge
                                                                  bcc.resize(ecnt); ecnt = 0;
class BCC_Bridge {
                                                                  for (int i = 0; i < n; ++i)
 private:
                                                                   if (not dfn[i]) dfs(i, i);
  int n, ecnt;
  vector<vector<pair<int,int>>> G;
                                                                 int get_id(int x) { return bcc[x]; }
  vector<int> dfn, low;
                                                                 int count() { return ecnt; }
  vector<bool> bridge;
                                                                 bool is_ap(int x) { return ap[x]; }
  void dfs(int u, int f) {
  dfn[u] = low[u] = dfn[f] + 1;
                                                              } bcc_ap;
                                                              3.4 2-SAT (SCC)
   for (auto [v, t]: G[u]) {
    if (v == f) continue;
                                                              class TwoSat{
    if (dfn[v]) {
                                                               private:
     low[u] = min(low[u], dfn[v]);
                                                                 int n;
     continue:
                                                                 vector<vector<int>> rG,G,sccs;
                                                                 vector<int> ord,idx;
    dfs(v, u);
                                                                 vector<bool> vis,result;
    low[u] = min(low[u], low[v]);
                                                                 void dfs(int u){
    if (low[v] > dfn[u]) bridge[t] = true;
                                                                  vis[u]=true
                                                                  for(int v:G[u])
                                                                  if(!vis[v]) dfs(v);
 public:
                                                                  ord.push_back(u);
  void init(int n_) {
   G.clear(); G.resize(n = n_);
                                                                 void rdfs(int u){
   low.assign(n, ecnt = 0);
                                                                 vis[u]=false;idx[u]=sccs.size()-1;
   dfn.assign(n, 0);
                                                                  sccs.back().push_back(u);
                                                                  for(int v:rG[u])
  void add_edge(int u, int v) {
                                                                   if(vis[v])rdfs(v);
   G[u].emplace_back(v, ecnt);
   G[v].emplace_back(u, ecnt++);
                                                                public:
                                                                 void init(int n_){
  void solve() {
                                                                 n=n_;G.clear();G.resize(n);
   bridge.assign(ecnt, false);
                                                                  rG.clear();rG.resize(n)
   for (int i = 0; i < n; ++i)
                                                                  sccs.clear();ord.clear()
    if (not dfn[i]) dfs(i, i);
                                                                  idx.resize(n);result.resize(n);
  bool is_bridge(int x) { return bridge[x]; }
                                                                 void add_edge(int u,int v){
} bcc_bridge;
                                                                 G[u].push_back(v);rG[v].push_back(u);
3.3 BCC Vertex
                                                                 void orr(int x,int y){
class BCC AP {
                                                                  if ((x^y)==1)return
 private:
                                                                  add_edge(x^1,y); add_edge(y^1,x);
  int n, ecnt;
  vector<vector<pair<int,int>>> G;
                                                                 bool solve(){
  vector<int> bcc, dfn, low, st;
                                                                 vis.clear();vis.resize(n);
  vector<bool> ap, ins;
void dfs(int u, int f)
                                                                  for(int i=0;i<n;++i)</pre>
                                                                  if(not vis[i])dfs(i);
   dfn[u] = low[u] = dfn[f] + 1;
                                                                  reverse(ord.begin(),ord.end());
   int ch = 0;
                                                                  for (int u:ord){
   for (auto [v, t]: G[u]) if (v != f) {
                                                                   if(!vis[u])continue;
    if (not ins[t]) {
                                                                   sccs.push_back(vector<int>());
     st.push_back(t);
                                                                   rdfs(u);
     ins[t] = true;
                                                                  for(int i=0;i<n;i+=2)</pre>
    if (dfn[v]) {
                                                                  if(idx[i]==idx[i+1])
     low[u] = min(low[u], dfn[v]);
                                                                    return false
     continue:
                                                                  vector<bool> c(sccs.size());
    } ++ch; dfs(v, u);
                                                                  for(size_t i=0;i<sccs.size();++i){</pre>
    low[u] = min(low[u], low[v]);
                                                                   for(size_t j=0;j<sccs[i].size();++j){
  result[sccs[i][j]]=c[i];</pre>
    if (low[v] >= dfn[u]) {
```

c[idx[sccs[i][j]^1]]=!c[i];

u = fa[ s ][ 0 ];

```
res.emplace_back( tl[ g ], tl[ u ] + 1 );
                                                                 while ( chain[ v ] != chain[ g ] ) {
   return true;
                                                                   int s = chain_st[ chain[ v ] ];
 bool get(int x){return result[x];}
                                                                  res.emplace_back( tl[ s ], tl[ v ] + 1 );
  inline int get_id(int x){return idx[x];}
                                                                  v = fa[ s ][ 0 ];
  inline int count(){return sccs.size();}
                                                                 res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
} sat2;
                                                                 return res;
3.5 Lowbit Decomposition
                                                                 /* res : list of intervals from u to v
                                                                  * ( note only nodes work, not edge )
class LowbitDecomp{
private:
                                                                   * vector< PII >& path = tree.get_path( u , v )
int time_, chain_, LOG_N;
vector< vector< int > > G, fa;
                                                                   * for( auto [ 1, r ] : path ) {
                                                                   * 0-base [ 1, r )
vector< int > tl, tr, chain, chain_st;
 // chain_: number of chain
 // tl, tr[ u ] : subtree interval in the seq. of u
// chain_st[ u ] : head of the chain contains u
                                                               } tree;
 // chian[ u ] : chain id of the chain u is on
void predfs( int u, int f ) {
                                                               3.6 MaxClique
 chain[ u ] = 0;
  for ( int v : G[ u ] ) {
                                                               // contain a self loop u to u, than u won't in clique
                                                               template < size_t MAXN >
  if ( v == f ) continue;
   predfs( v, u );
                                                               class MaxClique{
   if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
                                                               private:
    chain[ u ] = chain[ v ];
                                                                using bits = bitset< MAXN >;
                                                                bits popped, G[ MAXN ], ans
                                                                size_t deg[ MAXN ], deo[ MAXN ], n;
 if ( not chain[ u ] )
   chain[ u ] = chain_ ++;
                                                                void sort_by_degree() {
                                                                 popped.reset();
void dfschain( int u, int f ) {
                                                                 for ( size_t i = 0 ; i < n ; ++ i )</pre>
 fa[ u ][ 0 ] = f;
for ( int i = 1 ; i < LOG_N ; ++ i</pre>
                                                                    deg[ i ] = G[ i ].count();
                                                                 for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;</pre>
   fa[`u][i] = fa[fa[u][i-1]][i-1];
                                                                    for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )</pre>
  tl[ u ] = time_++;
  if ( not chain_st[ chain[ u ] ] )
                                                                        mi = deg[id = j]
  chain_st[ chain[ u ] ] = u;
  for ( int v : G[ u ] )
                                                                    popped[ deo[ i ] = id ] = 1;
  if ( v != f and chain[ v ] == chain[ u ] )
                                                                    for( size_t u = G[ i ]._Find_first() ;
  u < n ; u = G[ i ]._Find_next( u ) )</pre>
  dfschain( v, u );
for ( int v : G[ u ] )
                                                                      -- deg[ u ];
   if ( v != f and chain[ v ] != chain[ u ] )
    dfschain( v, u );
                                                                void BK( bits R, bits P, bits X ) {
 tr[ u ] = time_;
                                                                 if (R.count()+P.count() <= ans.count()) return;</pre>
bool anc( int u, int v ) {
                                                                 if ( not P.count() and not X.count() )
 return tl[ u ] <= tl[ v ] and tr[ v ] <= tr[ u ];</pre>
                                                                   if ( R.count() > ans.count() ) ans = R;
                                                                   return:
public:
int lca( int u, int v ) {
                                                                 /* greedily chosse max degree as pivot
  if ( anc( u, v ) ) return u;
                                                                 bits cur = P | X; size_t pivot = 0, sz = 0;
  for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
                                                                 for ( size_t u = cur._Find_first() ;
  if ( not anc( fa[ u ][ i ], v ) )
  u = fa[ u ][ i ];
                                                                   u < n ; u = cur._Find_next( u )</pre>
                                                                    if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
 return fa[ u ][ 0 ];
                                                                 cur = P & ( ~G[ pivot ] );
                                                                  */ // or simply choose first
                                                                 bits cur = P & (~G[ ( P | X )._Find_first() ]);
void init( int n ) {
 fa.assign( ++n, vector< int >( LOG_N ) );
for ( LOG_N = 0 ; ( 1 << LOG_N ) < n ; ++ LOG_N );</pre>
                                                                 for ( size_t u = cur._Find_first()
                                                                   u < n ; u = cur._Find_next( u ) ) {
 G.clear(); G.resize( n );
                                                                   if ( R[ u ] ) continue;
                                                                  R[u] = 1;
 tl.assign( n, 0 ); tr.assign( n, 0 );
                                                                   BK( R, P & G[ u ], X & G[ u ] );
 chain.assig( n, 0 ); chain_st.assign( n, 0 );
                                                                   R[u] = P[u] = 0, X[u] = 1;
void add_edge( int u , int v ) {
                                                                 }
  // 1-base
 G[ u ].push_back( v );
                                                               public:
 G[ v ].push_back( u );
                                                                void init( size_t n_ ) {
void decompose(){
                                                                 for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                  G[ i ].reset();
 chain_ = 1;
 predfs( 1, 1 );
                                                                 ans.reset();
  time_{-} = 0;
                                                                void add_edges( int u, bits S ) { G[ u ] = S; }
void add_edge( int u, int v ) {
  G[ u ][ v ] = G[ v ][ u ] = 1;
  dfschain(1,1);
PII get_subtree(int u) { return {tl[ u ],tr[ u ] }; }
vector< PII > get_path( int u , int v ){
 vector< PII > res;
                                                                int solve() {
                                                                 sort_by_degree(); // or simply iota( deo... )
  int g = lca( u, v );
 while ( chain[ u ] != chain[ g ] ) {
                                                                 for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                   deg[ i ] = G[ i ].count()
  int s = chain_st[ chain[ u ] ];
   res.emplace_back( tl[ s ], tl[ u ] + 1 );
                                                                 bits pob, nob = 0; pob.set();
```

for (size\_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>

```
for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
                                                              inline bool cmp(const int &i, const int &j) {
   size_t v = deo[ i ];
                                                              return dfn[i] < dfn[j];</pre>
   bits tmp; tmp[ v ] = 1;
   BK( tmp, pob \& G[v], nob \& G[v]);

pob[v] = 0, nob[v] = 1;
                                                              void build(int vectrices[], int k) {
                                                               static int stk[MAX_N];
                                                               sort(vectrices, vectrices + k, cmp);
  return static_cast< int >( ans.count() );
                                                               stk[sz++] = 0;
                                                               for (int i = 0; i < k; ++i) {
                                                                int u = vectrices[i], lca = LCA(u, stk[sz - 1]);
                                                                if (lca == stk[sz - 1]) stk[sz++] = u;
3.7 MaxCliqueDyn
                                                                 while (sz >= 2 && dep[stk[sz - 2]] >= dep[lca]) {
constexpr int kN = 150;
                                                                  addEdge(stk[sz - 2], stk[sz - 1]);
struct MaxClique { // Maximum Clique
                                                                  sz--;
bitset<kN> a[kN], cs[kN];
int ans, sol[kN], q, cur[kN], d[kN], n;
void init(int _n) {
                                                                 if (stk[sz - 1] != lca) {
                                                                  addEdge(lca, stk[--sz]);
 n = _n; for (int i = 0; i < n; i++) a[i].reset();</pre>
                                                                  stk[sz++] = lca, vectrices[cnt++] = lca;
 void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
                                                                 stk[sz++] = u;
 void csort(vector<int> &r, vector<int> &c)
  int mx = 1, km = max(ans - q + 1, 1), t = 0,
    m = int(r.size())
                                                               for (int i = 0; i < sz - 1; ++i)
  cs[1].reset(); cs[2].reset();
                                                                addEdge(stk[i], stk[i + 1]);
  for (int i = 0; i < m; i++) {
  int p = r[i], k = 1;
                                                              3.9 Centroid Decomposition
   while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
                                                             struct Centroid {
   cs[k][p] = 1;
                                                               vector<vector<int64_t>> Dist;
   if (k < km) r[t++] = p;
                                                               vector<int> Parent, Depth;
                                                               vector<int64_t> Sub, Sub2;
  c.resize(m);
                                                               vector<int> Sz, Sz2;
  if(t) c[t-1] = 0;
                                                               Centroid(vector<vector<pair<int, int>>> g) {
  for (int k = km; k <= mx; k++) {</pre>
                                                                int N = g.size();
   for (int p = int(cs[k]._Find_first());
                                                                vector<bool> Vis(N);
      p < kN; p = int(cs[k]._Find_next(p))) {
                                                                vector<int> sz(N), mx(N);
    r[t] = p; c[t++] = k;
                                                                vector<int> Path;
                                                                Dist.resize(N)
  }
                                                                Parent.resize(N):
                                                                Depth.resize(N)
 void dfs(vector<int> &r, vector<int> &c, int 1,
                                                                auto DfsSz = [&](auto dfs, int x) -> void {
  bitset<kN> mask) {
                                                                 Vis[x] = true; sz[x] = 1; mx[x] = 0;
  while (!r.empty()) {
                                                                 for (auto [u, w] : g[x]) {
   int p = r.back(); r.pop_back();
                                                                  if (Vis[u]) continue;
   mask[p] = 0;
                                                                  dfs(dfs, u)
   if (q + c.back() <= ans) return;</pre>
                                                                  sz[x] += sz[u];
   cur[q++] = p;
                                                                  mx[x] = max(mx[x], sz[u]);
   vector<int> nr, nc;
   bitset<kN> nmask = mask & a[p];
                                                                 Path.push_back(x);
   for (int i : r)
                                                                }:
    if (a[p][i]) nr.push_back(i);
                                                                auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
   if (!nr.empty()) {
                                                                 -> void {
    if (1 < 4) {
                                                                 Dist[x].push_back(D);Vis[x] = true;
     for (int i : nr)
                                                                 for (auto [u, w] : g[x]) {
      d[i] = int((a[i] & nmask).count());
                                                                  if (Vis[u]) continue;
     sort(nr.begin(), nr.end(),
                                                                  dfs(dfs, u, D + w);
      [&](int x, int y)
                                                                 }
       return d[x] > d[y];
                                                                }:
      });
                                                                auto Dfs = [&]
                                                                 (auto dfs, int x, int D = 0, int p = -1)->void {
    csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
                                                                 Path.clear(); DfsSz(DfsSz, x);
   } else if (q > ans) {
                                                                 int M = Path.size();
    ans = q; copy(cur, cur + q, sol);
                                                                 int C = -1;
                                                                 for (int u : Path) {
  if (max(M - sz[u], mx[u]) * 2 <= M) C = u;</pre>
   c.pop_back(); q--;
  }
                                                                  Vis[u] = false;
 int solve(bitset<kN> mask) { // vertex mask
                                                                 DfsDist(DfsDist, C);
  vector<int> r, c;
                                                                 for (int u : Path) Vis[u] = false;
  for (int i = 0; i < n; i++)
                                                                 Parent[C] = p; Vis[C] = true;
   if (mask[i]) r.push_back(i);
                                                                 Depth[C] = D;
  for (int i = 0; i < n; i++)
                                                                 for (auto [u, w] : g[C]) {
   d[i] = int((a[i] & mask).count());
                                                                  if (Vis[u]) continue;
  sort(r.begin(), r.end(),
                                                                  dfs(dfs, u, D + 1, C);
   [&](int i, int j) { return d[i] > d[j]; });
  csort(r, c);
  dfs(r, c, 1, mask);
                                                               Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
Sz.resize(N); Sz2.resize(N);
  return ans; // sol[0 ~ ans-1]
} graph;
                                                               void Mark(int v) {
                                                                int x = v, z = -1
                                                                for (int i = Depth[v]; i >= 0; --i) {
```

#### 3.8 Virtural Tree

```
Sub[x] += Dist[v][i]; Sz[x]++;
                                                                 cycle.PB(v);
   if (z != -1) {
                                                                vst[v]++;
    Sub2[z] += Dist[v][i];
                                                               }
    Sz2[z]++;
                                                               reverse(ALL(edgeID));
                                                               edgeID.resize(SZ(cycle));
   z = x; x = Parent[x];
                                                               return mmc;
  }
                                                             } mmc;
 int64_t Query(int v) {
                                                             3.12 Mo's Algorithm on Tree
  int64_t res = 0;
  int x = v, z = -1;
                                                             int q; vector< int > G[N];
  for (int i = Depth[v]; i >= 0; --i) {
                                                             struct Que{
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
                                                              int u, v, id;
   if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
                                                             } que[ N ];
                                                             int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
   z = x; x = Parent[x];
                                                             void dfs( int u, int f ) {
                                                              dfn[ u ] = dfn_++; int saved_rbp = stk_;
  return res;
                                                              for ( int v : G[ u ] ) {
                                                               if ( v == f ) continue;
                                                               dfs( v, u );
3.10 Tree Hashing
                                                               if ( stk_ - saved_rbp < SQRT_N ) continue;</pre>
                                                               for ( ++ block_ ; stk_ != saved_rbp ; )
  block_id[ stk[ -- stk_ ] ] = block_;
uint64_t hsah(int u, int f) {
 uint64_t r = 127;
 for (int v : G[ u ]) if (v != f) {
  uint64_t hh = hsah(v, u);
                                                              stk[stk_+ ++] = u;
  r=(r+(hh*hh)%1010101333)%1011820613;
                                                             bool inPath[ N ];
                                                             void Diff( int u ) {
 return r;
}
                                                              if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
                                                              else { /*add this edge*/ }
3.11 Minimum Mean Cycle
/* minimum mean cycle O(VE) */
                                                             void traverse( int& origin_u, int u ) {
struct MMC{
                                                              for ( int g = lca( origin_u, u )
#define FZ(n) memset((n),0,sizeof(n))
                                                               origin_u != g ; origin_u = parent_of[ origin_u ] )
#define E 101010
                                                                Diff( origin_u );
#define V 1021
                                                              for (int v = u; v != origin_u; v = parent_of[v])
                                                               Diff( v );
#define inf 1e9
 struct Edge { int v,u; double c; };
                                                              origin_u = u;
 int n, m, prv[V][V], prve[V][V], vst[V];
                                                             void solve() {
 Edge e[E];
                                                              dfs( 1, 1 );
while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
 vector<int> edgeID, cycle, rho;
 double d[V][V];
                                                              sort( que, que + q, [](const Que& x, const Que& y) {
 void init( int _n ) { n = _n; m = 0; }
 // WARNING: TYPE matters
                                                               return tie( block_id[ x.u ], dfn[ x.v ] )
 void add_edge( int vi , int ui , double ci )
{ e[ m ++ ] = { vi , ui , ci }; }
                                                                   < tie( block_id[ y.u ], dfn[ y.v ] );
                                                              int U = 1, V = 1;
 void bellman_ford() {
                                                              for ( int i = 0 ; i < q ; ++ i ) {
  for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {</pre>
                                                               pass( U, que[ i ].u );
   fill(d[i+1], d[i+1]+n, inf);
                                                               pass( V, que[ i ].v );
   for(int j=0; j<m; j++) {</pre>
                                                                // we could get our answer of que[ i ].id
    int v = e[j].v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
                                                             }
                                                             /*
     d[i+1][u] = d[i][v]+e[j].c;
                                                             Method 2:
     prv[i+1][u] = v;
                                                             dfs u:
     prve[i+1][u] = j;
                                                              push u
   }
                                                              iterate subtree
  }
                                                              push u
                                                             Let P = LCA(u, v), and St(u) <= St(v)
                                                             if (P == u) query[St(u), St(v)]
 double solve(){
  // returns inf if no cycle, mmc otherwise
                                                             else query[Ed(u), St(v)], query[St(P), St(P)]
  double mmc=inf;
  int st = -1;
                                                             3.13 Minimum Steiner Tree
  bellman_ford();
  for(int i=0; i<n; i++) {</pre>
                                                             // Minimum Steiner Tree
                                                             // 0(V 3^T + V^2 2^T)
   double avg=-inf;
   for(int k=0; k<n; k++) {</pre>
                                                             struct SteinerTree{
    if(d[n][i]<inf-eps)</pre>
                                                             #define V 33
     avg=max(avg,(d[n][i]-d[k][i])/(n-k));\\
                                                             #define T 8
    else avg=max(avg,inf);
                                                             #define INF 1023456789
                                                              int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
                                                              void init( int _n ){
                                                               FZ(vst);edgeID.clear();cycle.clear();rho.clear();
  for (int i=n; !vst[st]; st=prv[i--][st]) {
                                                                 dst[ i ][ j ] = INF;
   vst[st]++
   edgeID.PB(prve[i][st]);
                                                                 dst[ i ][ i ] = 0;
   rho.PB(st);
                                                               }
  while (vst[st] != 2) {
                                                              void add_edge( int ui , int vi , int wi ){
                                                               dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
   int v = rho.back(); rho.pop_back();
```

```
dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
void shortest_path(){
  for( int k = 0 ; k < n ; k ++ )
for( int i = 0 ; i < n ; i ++ )
    for( int j = 0 ; j < n ; j ++ )
dst[ i ][ j ] = min( dst[ i ][ j ],
    dst[ i ][ k ] + dst[ k ][ j ] );</pre>
 int solve( const vector<int>& ter ){
  int t = (int)ter.size();
  for( int i = 0 ; i < ( 1 << t ) ; i ++ )</pre>
  for( int j = 0 ; j < n ; j ++ )
dp[ i ][ j ] = INF;</pre>
  for( int i = 0 ; i < n ; i ++ )</pre>
  dp[ 0 ][ i ] = 0;
for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){
  if( msk == ( msk & (-msk) ) ){
    int who = _{-}lg( msk );
    for( int i = 0 ; i < n ; i ++ )
dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
    continue:
   for( int i = 0 ; i < n ; i ++ )</pre>
    for( int submsk = ( msk - 1 ) & msk ; submsk ;
       submsk = ( submsk - 1 ) & msk )
dp[ msk ][ i ] = min( dp[ msk ][ i ],
                 dp[ submsk ][ i ] +
                 dp[ msk ^ submsk ][ i ] );
   for( int i = 0 ; i < n ; i ++ ){
    tdst[ i ] = INF;
    for( int i = 0 ; i < n ; i ++ )</pre>
    dp[ msk ][ i ] = tdst[ i ];
  int ans = INF
  for( int i = 0 ; i < n ; i ++ )</pre>
   ans = min( ans , dp[ ( 1 << t ) - 1 ][ i ] );
  return ans;
}
} solver;
3.14
      Directed Minimum Spanning Tree
template <typename T> struct DMST {
```

```
T g[maxn][maxn], fw[maxn];
int n, fr[maxn];
bool vis[maxn], inc[maxn];
void clear() {
 for(int i = 0; i < maxn; ++i) {</pre>
  for(int j = 0; j < maxn; ++j) g[i][j] = inf;
  vis[i] = inc[i] = false;
 }
void addEdge(int u,int v,T w){g[u][v]=min(g[u][v],w);}
T operator()(int root, int _n) {
 n = n; T ans = 0;
 if (dfs(root) != n) return -1;
 while (true) {
  for(int i = 1;i <= n;++i) fw[i] = inf, fr[i] = i;</pre>
  for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
   for (int j = 1; j <= n; ++j) {
  if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
     fw[i] = g[j][i]; fr[i] = j;
  int x = -1;
  for(int i = 1;i <= n;++i)if(i != root && !inc[i]){</pre>
   int j = i, c = 0;
   while(j!=root && fr[j]!=i && c<=n) ++c, j=fr[j];</pre>
   if (j == root || c > n) continue;
   else { x = i; break; }
  if (!~x) {
   for (int i = 1; i <= n; ++i)
    if (i != root && !inc[i]) ans += fw[i];
   return ans;
  int y = x;
```

```
for (int i = 1; i <= n; ++i) vis[i] = false;
    do {
     ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true;
    } while (y != x);
    inc[x] = false;
    for (int k = 1; k <= n; ++k) if (vis[k]) {</pre>
     for (int j = 1; j <= n; ++j) if (!vis[j]) {
  if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
}
      if (g[j][k] < inf && g[j][k]-fw[k] < g[j][x])
g[j][x] = g[j][k] - fw[k];</pre>
  }
  return ans;
 int dfs(int now) {
  int r = 1; vis[now] = true;
for (int i = 1; i <= n; ++i)</pre>
   if (g[now][i] < inf && !vis[i]) r += dfs(i);</pre>
};
3.15
        Dominator Tree
```

```
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
 // vertices are numbered from 0 to n - 1
 fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
 fill(fa, fa + n, -1); fill(val, val + n, -1);
 fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
 fill(dom, dom + n, -1); tk = 0;
for (int i = 0; i < n; ++i) {
  g[i].clear(); r[i].clear(); rdom[i].clear();
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
 rev[dfn[x] = tk] = x;
 fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
 for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
  r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
 if (fa[x] == x) return c ? -1 : x;
 int p = find(fa[x], 1);
 if (p == -1) return c ? fa[x] : val[x];
 if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
 fa[x] = p;
 return c ? p : val[x];
vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
// p[i] = -2 if i is unreachable from s
 dfs(s);
 for (int i = tk - 1; i >= 0; --i) {
  for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
  if (i) rdom[sdom[i]].push_back(i);
  for (int &u : rdom[i]) {
   int p = find(u);
   if (sdom[p] == i) dom[u] = i;
   else dom[u] = p;
  if (i) merge(i, rp[i]);
 vector<int> p(n, -2); p[s] = -1;
 for (int i = 1; i < tk; ++i)
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
 for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
 return p;
}}
```

# 4 Matching & Flow

## 4.1 Kuhn Munkres

```
class KM {
```

```
private:
                                                               n=_n; walked.reset();
static constexpr lld INF = 1LL << 60;</pre>
                                                               for(int i=0;i<n;i++){</pre>
                                                                X[i].clear();Y[i].clear();
vector<lld> h1,hr,slk;
vector<int> fl,fr,pre,qu;
                                                                fX[i]=fY[i]=-1;
vector<vector<lld>> w;
vector<bool> v1,vr;
int n, ql, qr;
                                                              void add_edge(int x, int y){
                                                               X[x].push_back(y); Y[y].push_back(y);
bool check(int x) {
 if (v1[x] = true, f1[x] != -1)
  return vr[qu[qr++] = f1[x]] = true;
                                                              int solve(){
 while (x != -1) swap(x, fr[fl[x] = pre[x]]);
                                                               int cnt = 0;
                                                               for(int i=0;i<n;i++){</pre>
 return false;
}
                                                                walked.reset();
void bfs(int s) {
                                                                if(dfs(i)) cnt++;
 fill(slk.begin(), slk.end(), INF);
 fill(vl.begin(), vl.end(), false);
fill(vr.begin(), vr.end(), false);
                                                               // return how many pair matched
                                                               return cnt;
 ql = qr = 0;
 qu[qr++] = s;
                                                             };
  vr[s] = true;
                                                             4.3 General Graph Matching
 while (true) {
                                                             const int N = 514, E = (2e5) * 2;
  11d d;
  while (ql < qr) {</pre>
                                                             struct Graph{
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
                                                              int to[E],bro[E],head[N],e;
     if(!v1[x]&&slk[x]>=(d=h1[x]+hr[y]-w[x][y])){
                                                              int lnk[N], vis[N], stp, n;
      if (pre[x] = y, d) slk[x] = d;
                                                              void init( int _n ){
      else if (!check(x)) return;
                                                               stp = 0; e = 1; n = _n;
                                                               for( int i = 0 ; i <= n ; i ++ )</pre>
                                                                head[i] = lnk[i] = vis[i] = 0;
  d = INF;
                                                              void add_edge(int u,int v){
  for (int x = 0; x < n; ++x)
                                                               // 1-base
    if (!v1[x] && d > slk[x]) d = slk[x];
                                                               to[e]=v,bro[e]=head[u],head[u]=e++;
   for (int x = 0; x < n; ++x) {
                                                               to[e]=u,bro[e]=head[v],head[v]=e++;
   if (v1[x]) h1[x] += d;
    else slk[x] -= d;
                                                              bool dfs(int x){
   if (vr[x]) hr[x] -= d;
                                                               vis[x]=stp;
                                                               for(int i=head[x];i;i=bro[i]){
   for (int x = 0; x < n; ++x)
                                                                int v=to[i];
    if (!v1[x] && !s1k[x] && !check(x)) return;
                                                                if(!lnk[v]){
                                                                 lnk[x]=v, lnk[v]=x;
                                                                 return true
public:
                                                                }else if(vis[lnk[v]]<stp){</pre>
void init( int n_ ) {
                                                                 int w=lnk[v];
 n = n_; qu.resize(n);
                                                                 lnk[x]=v, lnk[v]=x, lnk[w]=0;
 fl.clear(); fl.resize(n, -1);
                                                                 if(dfs(w)) return true
 fr.clear(); fr.resize(n, -1);
                                                                 lnk[w]=v, lnk[v]=w, lnk[x]=0;
 hr.clear(); hr.resize(n); hl.resize(n);
 w.clear(); w.resize(n, vector<1ld>(n));
 slk.resize(n); pre.resize(n);
                                                               return false;
 vl.resize(n); vr.resize(n);
                                                              int solve(){
void set_edge( int u, int v, lld x ) {w[u][v] = x;}
                                                               int ans = 0;
11d solve() {
                                                               for(int i=1;i<=n;i++)</pre>
 for (int i = 0; i < n; ++i)
                                                                if(not lnk[i]){
  hl[i] = *max_element(w[i].begin(), w[i].end());
                                                                 stp++; ans += dfs(i);
 for (int i = 0; i < n; ++i) bfs(i);</pre>
 11d res = 0;
                                                               return ans;
 for (int i = 0; i < n; ++i) res += w[i][f1[i]];</pre>
                                                             } graph;
 return res;
}
                                                                   Minimum Weight Matching (Clique version)
} km;
                                                             struct Graph {
4.2 Bipartite Matching
                                                              // 0-base (Perfect Match)
class BipartiteMatching{
                                                              int n, edge[MXN][MXN];
                                                              int match[MXN], dis[MXN], onstk[MXN];
private:
vector<int> X[N], Y[N];
                                                              vector<int> stk;
int fX[N], fY[N], n;
                                                              void init(int _n) {
bitset<N> walked;
                                                               n = _n;
                                                               for (int i=0; i<n; i++)</pre>
bool dfs(int x){
                                                                for (int j=0; j<n; j++)</pre>
 for(auto i:X[x]){
  if(walked[i])continue;
                                                                 edge[i][j] = 0;
  walked[i]=1;
  if(fY[i]==-1||dfs(fY[i])){
                                                              void set_edge(int u, int v, int w) {
   fY[i]=x;fX[x]=i;
                                                               edge[u][v] = edge[v][u] = w;
    return 1;
  }
                                                              bool SPFA(int u){
 }
                                                               if (onstk[u]) return true;
 return 0;
                                                               stk.PB(u);
                                                               onstk[u] = 1;
public:
                                                               for (int v=0; v<n; v++){</pre>
                                                                if (u != v && match[u] != v && !onstk[v]){
void init(int _n){
```

```
int m = match[v];
    if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
     dis[m] = dis[u] - edge[v][m] + edge[u][v];
     onstk[v] = 1:
     stk.PB(v);
     if (SPFA(m)) return true;
     stk.pop_back();
     onstk[v] = 0;
  onstk[u] = 0;
  stk.pop_back();
  return false;
 int solve() {
  // find a match
  for (int i=0; i<n; i+=2){</pre>
   match[i] = i+1;
   match[i+1] = i;
  while (true){
   int found = 0;
   for (int i=0; i<n; i++)</pre>
    dis[i] = onstk[i] = 0;
   for (int i=0; i<n; i++){
    stk.clear()
    if (!onstk[i] && SPFA(i)){
     found = 1
     while (SZ(stk)>=2){
      int u = stk.back(); stk.pop_back();
int v = stk.back(); stk.pop_back();
      match[u] = v;
      match[v] = u;
   if (!found) break;
  int ret = 0;
  for (int i=0; i<n; i++)</pre>
   ret += edge[i][match[i]];
  return ret>>1;
} graph;
     Minimum Cost Circulation
struct Edge { int to, cap, rev, cost; };
vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
bool mark[kN];
int NegativeCycle(int n) {
memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
 int upd = -1:
 for (int i = 0; i <= n; ++i)
  for (int j = 0; j < n; ++j) {
   int idx = 0;
   for (auto &e : g[j]) {
    if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
     dist[e.to] = dist[j] + e.cost;
     pv[e.to] = j, ed[e.to] = idx;
     if (i == n) {
      upd = j;
      while(!mark[upd])mark[upd]=1,upd=pv[upd];
      return upd;
     }
    idx++;
```

}

return -1;

int Solve(int n) { int rt = -1, ans = 0;

while (!mark[rt]) {

while ((rt = NegativeCycle(n)) >= 0) { memset(mark, false, sizeof(mark));

cyc.emplace\_back(pv[rt], ed[rt]);

vector<pair<int, int>> cyc;

```
mark[rt] = true;
 rt = pv[rt];
reverse(cyc.begin(), cyc.end());
int cap = kInf;
for (auto &i : cyc)
 auto &e = g[i.first][i.second];
 cap = min(cap, e.cap);
for (auto &i : cyc) {
 auto &e = g[i.first][i.second];
  e.cap -= cap;
 g[e.to][e.rev].cap += cap;
 ans += e.cost * cap;
return ans;
```

#### 4.6 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.

  - 2. For each edge (x,y,l,u), connect  $x\to y$  with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v \to T$  with capacity -in(v).
    - To maximize, connect  $t\to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f 
      eq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the
    - maximum flow from s to t is the answer. – To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' 
      eq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e \, + \, f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipartite graph(X, Y)
  - 1. Redirect every edge:  $y \to x$  if  $(x,y) \in M$ ,  $x \to y$  otherwise
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited. 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Construct super source S and sink T
  - 2. For each edge (x, y, c), connect  $x \to y$  with (cost, cap) = (c, 1) if c>0, otherwise connect  $y\to x$  with (cost, cap)=(-c,1)
  - 3. For each edge with  $c < \mathbf{0}$ , sum these cost as K, then increase d(y)by 1, decrease  $d(\boldsymbol{x})$  by 1
  - 4. For each vertex v with d(v) > 0, connect  $S \to v$  with (cost, cap) =(0, d(v))
  - 5. For each vertex v with d(v) < 0, connect  $v \to T$  with (cost, cap) =(0, -d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let  ${\cal K}$  be the sum of all weights
  - 3. Connect source  $s \to v, v \in G$  with capacity K
  - 4. For each edge (u, v, w) in G, connect  $u \to v$  and  $v \to u$  with capacity
  - 5. For  $v~\in~G$  , connect it with sink  $v~\rightarrow~t$  with capacity K~+~2T~- $\left(\sum_{e \in E(v)} w(e)\right) - 2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight w(u, v).
  - 2. Connect v 
    ightharpoonup v' with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to  $\boldsymbol{v}$
  - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
  - 1. If  $p_v>0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$ . 2. Create edge (u,v) with capacity w with w being the cost of choosing
  - $\boldsymbol{u}$  without choosing  $\boldsymbol{v}.$
  - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity  $c_x$  and create edge (s,y) with ca-
- 2. Create edge (x,y) with capacity  $c_{xy}$ . 3. Create edge (x,y) and edge (x',y') with capacity  $c_{xyx'y'}$ .

```
4.7
      Dinic
template <typename Cap = int64_t>
class Dinic{
private:
 struct Edge{
  int to, rev;
  Cap cap;
 int n, st, ed;
 vector<vector<Edge>> G;
 vector<int> lv, idx;
 bool BFS(){
  fill(lv.begin(), lv.end(), -1);
  queue<int> bfs;
  bfs.push(st); lv[st] = 0;
  while(!bfs.empty()){
   int u = bfs.front(); bfs.pop();
   for(auto e: G[u]){
    if(e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
    bfs.push(e.to); lv[e.to] = lv[u] + 1;
  }
  return (lv[ed]!=-1);
 Cap DFS(int u, Cap f){
  if(u == ed) return f;
  Cap ret = 0;
  for(int &i = idx[u]; i < (int)G[u].size(); ++i){</pre>
   auto &e = G[u][i];
   if(e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
   Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
   G[e.to][e.rev].cap += nf;
   if(f == 0) return ret;
  if(ret == 0) lv[u] = -1;
  return ret;
public:
 void init(int n_, int st_, int ed_){
 n = n_{-}, st = st_{-}, ed = ed_{-};
  G.resize(n); lv.resize(n)
  fill(G.begin(), G.end(), vector<Edge>());
 void add_edge(int u, int v, Cap c){
  G[u].push_back({v, (int)G[v].size(), c});
  G[v].push_back({u, ((int)G[u].size())-1, 0});
```

#### Minimum Cost Maximum Flow

Cap f = DFS(st, numeric\_limits<Cap>::max());

Cap max\_flow(){

Cap ret = 0;

ret += f;

return ret:

}

};

while(BFS()){

idx.assign(n, 0);

if(f == 0) break;

```
class MiniCostMaxiFlow{
using Cap = int; using Wei = int64_t;
using PCW = pair<Cap,Wei>;
static constexpr Cap INF_CAP = 1 << 30;</pre>
static constexpr Wei INF_WEI = 1LL<<60;</pre>
private:
struct Edge{
 int to, back;
 Cap cap; Wei wei;
 Edge() {}
 Edge(int a,int b, Cap c, Wei d):
   to(a),back(b),cap(c),wei(d)
  {}
 };
int ori, edd;
vector<vector<Edge>> G;
vector<int> fa, wh;
vector<bool> inq;
vector<Wei> dis;
PCW SPFA(){
```

```
fill(inq.begin(),inq.end(),false);
  fill(dis.begin(), dis.end(), INF_WEI);
  queue<int> qq; qq.push(ori);
  dis[ori]=0;
  while(!qq.empty()){
   int u=qq.front();qq.pop();
   inq[u] = 0;
   for(int i=0;i<SZ(G[u]);++i){</pre>
    Edge e=G[u][i];
    int v=e.to;
    Wei d=e.wei;
    if(e.cap<=0||dis[v]<=dis[u]+d)
     continue:
    dis[v]=dis[u]+d;
    fa[v]=u,wh[v]=i;
    if(inq[v]) continue;
    qq.push(v);
    inq[v]=1;
   }
  if(dis[edd]==INF_WEI) return {-1, -1};
  Cap mw=INF_CAP;
  for(int i=edd;i!=ori;i=fa[i])
   mw=min(mw,G[fa[i]][wh[i]].cap);
  for (int i=edd;i!=ori;i=fa[i]){
   auto &eg=G[fa[i]][wh[i]];
   eq.cap-=mw;
   G[eg.to][eg.back].cap+=mw;
  }
  return {mw,dis[edd]};
public:
 void init(int a,int b,int n){
  ori=a,edd=b;
  G.clear();G.resize(n);
  fa.resize(n);wh.resize(n);
  inq.resize(n); dis.resize(n);
 void add_edge(int st, int ed, Cap c, Wei w){
  G[st].emplace_back(ed,SZ(G[ed]),c,w);
  G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
 PCW solve(){
  /* might modify to
  cc += ret.first * ret.second
  ww += ret.first * ret.second
  */
  Cap cc=0; Wei ww=0;
  while(true){
   PCW ret=SPFA();
   if(ret.first==-1) break;
   cc+=ret.first;
   ww+=ret.second;
  return {cc,ww};
 }
} mcmf;
4.9 Global Min-Cut
```

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
}
pair<int, int> phase(int n) {
memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
 int s = -1, t = -1;
 while (true) {
  int c = -1;
  for (int i = 0; i < n; ++i) {
   if (del[i] || v[i]) continue;
   if (c == -1 \mid | g[i] > g[c]) c = i;
  if (c == -1) break;
  v[s = t, t = c] = true;
  for (int i = 0; i < n; ++i) {
   if (del[i] || v[i]) continue;
   g[i] += w[c][i];
```

return s-1;

```
void init(){
 return make_pair(s, t);
                                                                     primes.reserve(N);
                                                                     primes.push_back(1);
                                                                     for(int i=2;i<N;i++) {</pre>
int mincut(int n) {
int cut = 1e9:
                                                                      if(!sieved[i]) primes.push_back(i);
 memset(del, false, sizeof(del));
                                                                      pi[i] = !sieved[i] + pi[i-1];
 for (int i = 0; i < n - 1; ++i) {
                                                                      for(int p: primes) if(p > 1) {
  int s, t; tie(s, t) = phase(n);
                                                                       if(p * i >= N) break;
                                                                       sieved[p * i] = true;
  del[t] = true; cut = min(cut, g[t]);
  for (int j = 0; j < n; ++j) {
  w[s][j] += w[t][j]; w[j][s] += w[j][t];
                                                                       if(p % i == 0) break;
  }
 }
 return cut;
                                                                    11d phi(11d m, 11d n) {
                                                                     static constexpr int MM = 80000, NN = 500;
                                                                     static lld val[MM][NN];
5
     Math
                                                                     if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
                                                                     if(n == 0) return m;
      Prime Table
                                                                     if(primes[n] >= m) return 1;
1002939109, 1020288887, 1028798297, 1038684299,
                                                                     1ld ret = phi(m,n-1)-phi(m/primes[n],n-1);
1041211027, 1051762951, 1058585963, 1063020809,\\
                                                                     if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
1147930723, 1172520109, 1183835981, 1187659051,
                                                                     return ret;
1241251303, 1247184097, 1255940849, 1272759031,\\
1287027493, 1288511629, 1294632499, 1312650799,\\
1868732623, 1884198443, 1884616807, 1885059541, 1909942399, 1914471137, 1923951707, 1925453197,
                                                                    11d pi_count(11d);
                                                                    11d P2(11d m, 11d n) {
1979612177, 1980446837, 1989761941, 2007826547,
                                                                     11d sm = square_root(m), ret = 0;
2008033571, 2011186739, 2039465081, 2039728567,\\
                                                                     for(lld i = n+1;primes[i]<=sm;i++)</pre>
2093735719, 2116097521, 2123852629, 2140170259\\
                                                                      ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
\begin{array}{c} 3148478261, 3153064147, 3176351071, 3187523093, \\ 3196772239, 3201312913, 3203063977, 3204840059, \end{array}
                                                                     return ret:
3210224309, 3213032591, 3217689851, 3218469083, \\ 3219857533, 3231880427, 3235951699, 3273767923, \\
                                                                    11d pi_count(11d m) {
3276188869, 3277183181, 3282463507, 3285553889,
                                                                     if(m < N) return pi[m];</pre>
3319309027, 3327005333, 3327574903, 3341387953,
                                                                     11d n = pi_count(cube_root(m));
3373293941, 3380077549, 3380892997, 3381118801\\
                                                                     return phi(m, n) + n - 1 - P2(m, n);
     \lfloor rac{n}{i} 
floor Enumeration
T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor
                                                                    5.6 Range Sieve
5.3 ax+by=gcd
// ax+ny = 1, ax+ny == ax == 1 \pmod{n}
                                                                    const int MAX_SQRT_B = 50000;
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
                                                                    const int MAX_L = 200000 + 5;
 if (y == 0) g=x,a=1,b=0;
 else exgcd(y,x%y,g,b,a),b=(x/y)*a;
                                                                    bool is_prime_small[MAX_SQRT_B];
                                                                    bool is_prime[MAX_L];
5.4 Pollard Rho
                                                                    void sieve(lld 1, lld r){
// does not work when n is prime
                                                                     // [1, r)
// return any non-trivial factor
                                                                     for(lld i=2;i*i<r;i++) is_prime_small[i] = true;</pre>
                                                                     for(lld i=1;i<r;i++) is_prime[i-1] = true;</pre>
llu pollard_rho(llu n){
                                                                     if(l==1) is_prime[0] = false;
 static auto f=[](llu x,llu k,llu m){
  return add(k,mul(x,x,m),m);
                                                                     for(lld i=2;i*i<r;i++){</pre>
                                                                      if(!is_prime_small[i]) continue;
 }:
 if (!(n&1)) return 2;
                                                                      for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;</pre>
 mt19937 rnd(120821011);
                                                                      for(1ld j=std::max(2LL, (1+i-1)/i)*i;j<r;j+=i)</pre>
                                                                         is_prime[j-1]=false;
 while(true){
  llu y=2,yy=y,x=rnd()%n,t=1;
  for(llu sz=2;t==1;sz<<=1) {</pre>
                                                                   }
   for(llu i=0;i<sz;++i){</pre>
    if(t!=1)break;
                                                                    5.7 Miller Rabin
    yy=f(yy,x,n);
                                                                    bool isprime(llu x){
    t=gcd(yy>y?yy-y:y-yy,n);
                                                                     static llu magic[]={2,325,9375,28178,\
                                                                                450775,9780504,1795265022};
   y=yy;
                                                                     static auto witn=[](llu a,llu u,llu n,int t)
                                                                     ->bool{
  if(t!=1&&t!=n) return t;
                                                                      if (!(a = mpow(a%n,u,n)))return 0;
                                                                      while(t--)
                                                                       llu a2=mul(a,a,n);
     Pi Count (Linear Sieve)
                                                                       if(a2==1 && a!=1 && a!=n-1)
                                                                        return 1;
static constexpr int N = 1000000 + 5;
                                                                       a = a2;
lld pi[N];
                                                                      }
vector<int> primes;
                                                                      return a!=1;
bool sieved[N];
11d cube_root(11d x){
                                                                     if(x<2)return 0;</pre>
 1ld s=cbrt(x-static_cast<long double>(0.1));
                                                                     if(!(x&1))return x==2;
 while(s*s*s <= x) ++s;
                                                                     llu x1=x-1; int t=0;
 return s-1;
                                                                     while(!(x1&1))x1>>=1,t++;
                                                                     for(llu m:magic)if(witn(m,x1,x,t))return 0;
11d square_root(11d x){
                                                                     return 1;
1ld s=sqrt(x-static_cast<long double>(0.1));
 while(s*s \ll x) ++s;
```

#### 5.8 Inverse Element

```
// x's inverse mod k
                                                                 return cplx(cos(2*pi*j/maxn), sin(2*pi*j/maxn));
long long GetInv(long long x, long long k){
                                                                });
// k is prime: euler_(k)=k-1
                                                               }
                                                               void fft(vector<cplx> &v, int n) {
 return qPow(x, euler_phi(k)-1);
                                                               int z = __builtin_ctz(n) - 1;
// if you need [1, x] (most use: [1, k-1]
void solve(int x, long long k){
                                                                for (int i = 0; i < n; ++i) {
                                                                 int x = 0, j = 0;
                                                                 for (;(1 << j) < n;++j) x^=(i >> j & 1)<<(z - j);
inv[1] = 1;
 for(int i=2;i<x;i++)</pre>
                                                                 if (x > i) swap(v[x], v[i]);
  inv[i] = ((long long)(k - k/i) * inv[k % i]) % k;
                                                                for (int s = 2; s <= n; s <<= 1) {
                                                                 int z = s >> 1;
5.9 Euler Phi Function
                                                                 for (int i = 0; i < n; i += s) {</pre>
                                                                  for (int k = 0; k < z; ++k) {
                                                                   cplx x = v[i + z + k] * omega[maxn / s * k];
  extended euler:
                                                                   v[i + z + k] = v[i + k] - x;
  a^b mod p
                                                                   v[i + k] = v[i + k] + x;
  if gcd(a, p)==1: a^{(b\%phi(p))}
  elif b < phi(p): a^b mod p
  else a^(b%phi(p) + phi(p))
lld euler_phi(int x){
                                                               void ifft(vector<cplx> &v, int n) {
 11d r=1;
                                                                fft(v, n);
 for(int i=2;i*i<=x;++i){</pre>
                                                               reverse(v.begin() + 1, v.end());
  if(x\%i==0){
                                                                for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);
   x/=i; r*=(i-1);
   while(x%i==0){
                                                              VL convolution(const VI &a, const VI &b) {
    x/=i; r*=i;
                                                                // Should be able to handle N <= 10^5, C <= 10^4
                                                                int sz = 1:
  }
                                                                while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
                                                                vector<cplx> v(sz);
 if(x>1) r*=x-1;
                                                                for (int i = 0; i < sz; ++i) {
 return r;
                                                                 double re = i < a.size() ? a[i] : 0;</pre>
                                                                 double im = i < b.size() ? b[i] : 0;</pre>
vector<int> primes;
                                                                 v[i] = cplx(re, im);
bool notprime[N];
1ld phi[N];
                                                                fft(v, sz);
void euler_sieve(int n){
                                                                for (int i = 0; i <= sz / 2; ++i) {
 for(int i=2;i<n;i++){</pre>
                                                                 int j = (sz - i) & (sz - 1);
  if(!notprime[i]){
                                                                 cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
   primes.push_back(i); phi[i] = i-1;
                                                                   * cplx(0, -0.25);
                                                                 if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
  for(auto j: primes){
                                                                   ].conj()) * cplx(\hat{0}, -0.25);
   if(i*j >= n) break;
                                                                 v[i] = x;
   notprime[i*j] = true;
   phi[i*j] = phi[i] * phi[j];
                                                                ifft(v, sz);
   if(i % j == 0){
                                                                VL c(sz);
    phi[i*j] = phi[i] * j;
                                                                for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
    break:
                                                               VI convolution_mod(const VI &a, const VI &b, int p) {
                                                                int sz = 1;
                                                                while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
                                                                vector<cplx> fa(sz), fb(sz);
for (int i = 0; i < (int)a.size(); ++i)</pre>
5.10
      Gauss Elimination
void gauss(vector<vector<double>> &d) {
                                                                fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 int n = d.size(), m = d[0].size();
                                                                for (int i = 0; i < (int)b.size(); ++i)</pre>
 for (int i = 0; i < m; ++i) {
                                                                 fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
  int p = -1;
                                                                fft(fa, sz), fft(fb, sz);
  for (int j = i; j < n; ++j) {
  if (fabs(d[j][i]) < eps) continue;</pre>
                                                                double r = 0.25 / sz;
                                                                cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
   if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
                                                                 int j = (sz - i) & (sz - 1);
  if (p == -1) continue;
                                                                 cplx a1 = (fa[i] + fa[j].conj());
  for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
                                                                 cplx a2 = (fa[i] - fa[j].conj()) * r2;
  for (int j = 0; j < n; ++j) {
                                                                 cplx b1 = (fb[i] + fb[j].conj()) * r3;
  if (i == j) continue;
                                                                 cplx b2 = (fb[i] - fb[j].conj()) * r4;
   double z = d[j][i] / d[i][i];
                                                                 if (i != j)
   for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
                                                                  cplx c1 = (fa[j] + fa[i].conj());
cplx c2 = (fa[j] - fa[i].conj()) * r2;
                                                                  cplx d1 = (fb[j] + fb[i].conj()) * r3;
                                                                  cplx d2 = (fb[j] - fb[i].conj()) * r4;
                                                                  fa[i] = c1 * d1 + c2 * d2 * r5;
5.11
      Fast Fourier Transform
                                                                  fb[i] = c1 * d2 + c2 * d1;
namespace fft {
using VI = vector<int>;
                                                                 fa[j] = a1 * b1 + a2 * b2 * r5;
                                                                 fb[j] = a1 * b2 + a2 * b1;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
                                                                fft(fa, sz), fft(fb, sz);
                                                                vector<int> res(sz);
void prefft() {
 generate_n(omega, maxn + 1, [i=0]()mutable{
                                                                for (int i = 0; i < sz; ++i) {
 auto j = i++;
                                                                long long a = round(fa[i].re), b = round(fb[i].re),
```

```
c = round(fa[i].im);
                                                                 if (inv) {
  res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
                                                                  int invn = modinv(n);
                                                                  for (int i = 0; i < n; i++)</pre>
                                                                   F[i] = modmul(F[i], invn);
 return res;
}}
                                                                  reverse(F + 1, F + n);
                                                                }
5.12 Chinese Remainder
                                                              };
1ld crt(lld ans[], lld pri[], int n){
                                                              const int P=2013265921, root=31;
 11d M = 1, ret = 0;
                                                              const int MAXN=1<<20;</pre>
 for(int i=0;i<n;i++) M *= pri[i];</pre>
                                                              NTT<P, root, MAXN> ntt;
 for(int i=0;i<n;i++){</pre>
  lld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
                                                              5.15 Polynomial Operations
  ret += (ans[i]*(M/pri[i])%M * iv)%M;
                                                              using VL = vector<LL>;
  ret %= M;
                                                              #define fi(s, n) for (int i=int(s); i<int(n); ++i)</pre>
                                                              #define Fi(s, n) for (int i=int(n); i>int(s); --i)
 return ret;
                                                              int n2k(int n) {
                                                               int sz = 1; while (sz < n) sz <<= 1;</pre>
/*
                                                                return sz;
Another:
x = a1 \% m1
                                                              template<int MAXN, LL P, LL RT> // MAXN = 2^k
x = a2 \% m2
                                                              struct Poly { // coefficients in [0, P)
g = gcd(m1, m2)
                                                                static NTT<MAXN, P, RT> ntt;
assert((a1-a2)%g==0)
                                                               VL coef:
[p, q] = exgcd(m2/g, m1/g)
                                                                int n() const { return coef.size(); } // n()>=1
return a2+m2*(p*(a1-a2)/g)
                                                                LL *data() { return coef.data(); }
0 <= x < lcm(m1, m2)
                                                                const LL *data() const { return coef.data(); }
*/
                                                                LL &operator[](size_t i) { return coef[i]; }
                                                                const LL &operator[](size_t i)const{return coef[i];}
      Berlekamp Massey
                                                                Poly(initializer_list<LL> a) : coef(a) { }
// x: 1-base, p[]: 0-base
                                                                explicit Poly(int _n = 1) : coef(_n) { }
template<size_t N>
                                                               Poly(const LL *arr, int _n) : coef(arr, arr + _n) {}
Poly(const Poly &p, int _n) : coef(_n) {
vector<llf> BM(llf x[N], size_t n){
 size_t f[N]={0},t=0;11f d[N];
                                                                copy_n(p.data(), min(p.n(), _n), data());
 vector<llf> p[N];
 for(size_t i=1,b=0;i<=n;++i) {</pre>
                                                                Poly& irev(){return reverse(data(),data()+n()),*this;}
 for(size_t j=0;j<p[t].size();++j)
d[i]+=x[i-j-1]*p[t][j];</pre>
                                                                Poly& isz(int _n) { return coef.resize(_n), *this; }
                                                                Poly& iadd(const Poly &rhs) { // n() == rhs.n()
  if(abs(d[i]-=x[i])<=EPS)continue;</pre>
                                                                fi(0, n()) if ((coef[i]+=rhs[i]) >= P)coef[i]-=P;
  f[t]=i; if(!t){p[++t].resize(i); continue;}
                                                                 return *this;
  vector<llf> cur(i-f[b]-1);
  llf k=-d[i]/d[f[b]]; cur.PB(-k);
                                                                Poly& imul(LL k) {
  for(size_t j=0;j<p[b].size();j++)</pre>
                                                                fi(0, n()) coef[i] = coef[i] * k % P;
   cur.PB(p[b][j]*k);
                                                                 return *this;
  if(cur.size()<p[t].size())cur.resize(p[t].size());</pre>
  for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];</pre>
                                                                Poly Mul(const Poly &rhs) const {
  if(i-f[b]+p[b].size()>=p[t].size()) b=t;
                                                                const int _n = n2k(n() + rhs.n() - 1);
                                                                Poly X(*this, _n), Y(rhs, _n);
ntt(X.data(), _n), ntt(Y.data(),
fi(0, _n) X[i] = X[i] * Y[i] % P;
 p[++t]=cur;
return p[t];
                                                                 ntt(X.data(), _n, true);
                                                                 return X.isz(n() + rhs.n() - 1);
5.14 NTT
template <int mod, int G, int maxn>
                                                                Poly Inv() const { // coef[0] != 0
                                                                if (n() == 1) return {ntt.minv(coef[0])};
const int _n = n2k(n() * 2);
struct NTT {
 static_assert (maxn == (maxn & -maxn));
                                                                 Poly Xi = Poly(*this, (n() + 1)/2).Inv().isz(_n);
 int roots[maxn];
                                                                 Poly Y(*this, _n);
 NTT () {
  int r = modpow(G, (mod - 1) / maxn);
                                                                 ntt(Xi.data(), _n), ntt(Y.data(), _n);
                                                                 fi(0, _n) {
Xi[i] *= (2 - Xi[i] * Y[i]) % P
  for (int i = maxn >> 1; i; i >>= 1) {
   roots[i] = 1;
   for (int j = 1; j < i; j++)
                                                                  if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
    roots[i + j] = modmul(roots[i + j - 1], r);
   r = modmul(r, r);
                                                                ntt(Xi.data(), _n, true);
                                                                 return Xi.isz(n());
 }
Poly Sqrt() const { // Jacobi(coef[0], P) = 1
                                                                if (n()==1) return {QuadraticResidue(coef[0], P)};
Poly X = Poly(*this, (n()+1) / 2).Sqrt().isz(n());
   if (i < j) swap(F[i], F[j]);</pre>
                                                                 return X.iadd(Mul(X.Inv()).isz(n())).imul(P/2+1);
   for (int k = n > 1; (j^k < k; k > = 1);
                                                                pair<Poly, Poly> DivMod(const Poly &rhs) const {
  for (int s = 1; s < n; s *= 2) {
                                                                 // (rhs.)back() != 0
   for (int i = 0; i < n; i += s * 2) {
                                                                 if (n() < rhs.n()) return {{0}, *this};</pre>
    for (int j = 0; j < s; j++) {
                                                                 const int _n = n() - rhs.n() + 1;
     int a = F[i+j];
                                                                 Poly X(rhs); X.irev().isz(_n);
     int b = modmul(F[i+j+s], roots[s+j]);
                                                                 Poly Y(*this); Y.irev().isz(_n);
                                                                 Poly Q = Y.Mul(X.Inv()).isz(_n).irev();
     F[i+j] = modadd(a, b); // a + b
     F[i+j+s] = modsub(a, b); // a - b
                                                                 X = rhs.Mul(Q), Y = *this
                                                                 fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
                                                                 return {Q, Y.isz(max(1, rhs.n() - 1))};
```

```
Poly Dx() const {
                                                              return R.isz(L);
Poly ret(n() - 1);
 fi(0, ret.n()) ret[i] = (i + 1) * coef[i + 1] % P;
                                                             static LL LinearRecursion(const VL&a,const VL&c,LL n){
 return ret.isz(max(1, ret.n()));
                                                              // a_n = \sum_{j=0}^{n-j} a_{n-j}
                                                              const int k = (int)a.size();
Poly Sx() const {
                                                              assert((int)c.size() == k + 1);
Poly ret(n() + 1);
                                                              Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
fi(0, n()) ret[i + 1]=ntt.minv(i + 1)*coef[i] % P;
                                                              fi(1, k + 1) C[k - i] = c[i] ? P - c[i] : 0;
 return ret;
                                                              C[k] = 1
                                                              while (n) {
                                                               if (n % 2) W = W.Mul(M).DivMod(C).second;
Poly _tmul(int nn, const Poly &rhs) const {
Poly Y = Mul(rhs).isz(n() + nn - 1);
                                                               n /= 2, M = M.Mul(M).DivMod(C).second;
return Poly(Y.data() + n() - 1, nn);
                                                              LL ret = 0;
                                                              fi(0, k) ret = (ret + W[i] * a[i]) % P;
VL _eval(const VL &x, const auto up)const{
const int _n = (int)x.size();
                                                              return ret;
 if (!_n) return {};
vector<Poly> down(_n * 2);
                                                            }:
 down[1] = DivMod(up[1]).second;
                                                            #undef fi
 fi(2,_n*2) down[i]=down[i/2].DivMod(up[i]).second;
                                                            #undef Fi
                                                            using Poly_t = Poly<131072 * 2, 998244353, 3>;
 /* down[1] = Poly(up[1]).irev().isz(n()).Inv().irev()
     _tmul(_n, *this)
                                                            template<> decltype(Poly_t::ntt) Poly_t::ntt = {};
 fi(2, _n * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
                                                            5.16 FWT
   1, down[i / 2]); */
VL y(_n);
                                                            /* xor convolution:
fi(0, _n) y[i] = down[_n + i][0];
                                                             * x = (x0, x1) , y = (y0, y1)
 return y;
                                                             *z = (x0y0 + x1y1 , x0y1 + x1y0 )
static vector<Poly> _tree1(const VL &x) {
                                                             * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
const int _n = (int)x.size();
                                                             * z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
 vector<Poly> up(_n * 2);
                                                             *z = (1/2) *z'
 fi(0, _n) up[_n + i] = \{(x[i] ? P - x[i] : 0), 1\};
                                                             * or convolution:
Fi(0, _n-1) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
                                                             * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
return up;
                                                             * and convolution:
                                                             * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
VL Eval(const VL&x)const{return _eval(x,_tree1(x));}
                                                            const LL MOD = 1e9+7;
static Poly Interpolate(const VL &x, const VL &y) {
                                                            inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
 const int _n = (int)x.size();
                                                             for( int d = 1 ; d < N ; d <<= 1 ) {
vector<Poly> up = _tree1(x), down(_n * 2);
VL z = up[1].Dx()._eval(x, up);
                                                              int d2 = d << 1;
                                                              for( int s = 0 ; s < N ; s += d2 )
 fi(0, _n) z[i] = y[i] * ntt.minv(z[i]) % P;
                                                               for( int i = s , j = s+d ; i < s+d ; i++, j++ ){ LL ta = x[i] , tb = x[j];
fi(0, n) down[n+1] = \{z[i]\};

Fi(0, n-1) down[i] = down[i * 2].Mul(up[i * 2 + 1])
                                                                x[ i ] = ta+tb;
  .iadd(down[i * 2 + 1].Mul(up[i * 2]));
                                                                x[ j ] = ta-tb;
if( x[ i ] >= MOD ) x[ i ] -= MOD;
 return down[1];
                                                                if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
Poly Ln() const { // coef[0] == 1
return Dx().Mul(Inv()).Sx().isz(n());
                                                             if( inv )
Poly Exp() const \{ // coef[0] == 0 \}
                                                              for( int i = 0 ; i < N ; i++ ) {
  x[ i ] *= inv( N, MOD );</pre>
if (n() == 1) return {1};
 Poly X = Poly(*this, (n() + 1)/2).Exp().isz(n());
                                                               x[ i ] %= MOD;
Poly Y = X.Ln(); Y[0] = P - 1;
                                                              }
fi(0, n()) if((Y[i] = coef[i] - Y[i]) < 0)Y[i]+=P;
 return X.Mul(Y).isz(n());
                                                                   DiscreteLog
                                                            5.17
Poly Pow(const string &K) const {
                                                            11d BSGS(11d P, 11d B, 11d N) {
int nz = 0;
                                                             // find B^L = N mod P
 while (nz < n() && !coef[nz]) ++nz;</pre>
                                                             unordered_map<lld, int> R;
 LL nk = 0, nk2 = 0;
                                                             11d sq = (11d)sqrt(P);
 for (char c : K) {
                                                             11d t = 1;
 nk = (nk * 10 + c - '0') % P;
                                                             for (int i = 0; i < sq; i++) {
  nk2 = nk2 * 10 + c - '0';
                                                              if (t == N) return i;
 if (nk2 * nz >= n()) return Poly(n());
                                                              if (!R.count(t)) R[t] = i;
  nk2 %= P - 1;
                                                              t = (t * B) % P;
 if (!nk && !nk2) return Poly({1}, n());
                                                             11d f = inverse(t, P);
Poly X(data() + nz, n() - nz * nk2);
                                                             for(int i=0;i<=sq+1;i++) {</pre>
LL \times 0 = X[0];
                                                              if (R.count(N))
 return X.imul(ntt.minv(x0)).Ln().imul(nk).Exp()
                                                               return i * sq + R[N];
  .imul(ntt.mpow(x0, nk2)).irev().isz(n()).irev();
                                                              N = (N * f) % P;
Poly InvMod(int L) { // (to evaluate linear recursion)
                                                             return -1;
Poly R{1, 0}; // *this * R mod x^L = 1 (*this[0] ==
 for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                            5.18 FloorSum
  Poly 0 = R.Mul(Poly(data(), min(2 << level, n())));
                                                            // @param n `n < 2^32`
  Poly Q(2 << level); Q[0] = 1;
                                                            // @param m `1 <= m < 2^32`
  for (int j = (1 << level); j < (2 << level); ++j)</pre>
  Q[j] = (P - O[j]) \% P;
                                                            // @return sum_\{i=0\}^{n-1} floor((ai + b)/m) mod 2^64
  R = R.Mul(Q).isz(4 << level);
                                                            llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
                                                            11u ans = 0;
```

aux[t] = i;

```
while (true)
                                                                        db(t + 1, t, n, k);
  if (a >= m)
   ans += n * (n - 1) / 2 * (a / m); a %= m;
                                                                     }
  if (b >= m) {
                                                                    int de_bruijn(int k, int n) {
                                                                     // return cyclic string of len k^n s.t. every string
   ans += n * (b / m); b %= m;
                                                                      // of len n using k char appears as a substring.
                                                                     if (k == 1) {
  11u y_max = a * n + b;
  if (y_max < m) break;</pre>
                                                                      res[0] = 0;
 // y_max < m * (n + 1)
// floor(y_max / m) <= n
                                                                      return 1:
  n = (1lu)(y_max / m), b = (1lu)(y_max % m);
                                                                     for (int i = 0; i < k * n; i++) aux[i] = 0;
                                                                      sz = 0;
  swap(m, a);
                                                                      db(1, 1, n, k);
                                                                      return sz;
 return ans;
11d floor_sum(11d n, 11d m, 11d a, 11d b) {
                                                                    5.21 Simplex Construction
 assert(0 <= n && n < (1LL << 32));
 assert(1 <= m && m < (1LL << 32));
                                                                    Standard form: maximize \sum_{1 \le i \le n} c_i x_i such that for all 1 \le j \le m,
 llu ans = 0;
                                                                    \sum_{1 \le i \le n} A_{ji} x_i \le b_j and x_i \ge 0 for all 1 \le i \le n.
 if (a < 0) {
  11u a2 = (a \% m + m) \% m;
                                                                       1. In case of minimization, let c_i' = -c_i
  ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
                                                                       2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
  a = a2:
                                                                       3. \sum_{1 \leq i \leq n} A_{ji} x_i = b_j
 if (b < 0) {
 11\dot{u} b2 = (b % m + m) % m;
                                                                             • \sum_{1 \le i \le n} A_{ji} x_i \le b_j
  ans -= 1ULL * n * ((b2 - b) / m);
                                                                             • \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j
  b = b2;
                                                                       4. If x_i has no lower bound, replace x_i with x_i - x_i'
return ans + floor_sum_unsigned(n, m, a, b);
                                                                    5.22 Simplex
5.19 Quadratic residue
                                                                    namespace simplex {
struct Status{
                                                                    // maximize c^Tx under Ax <= B
                                                                    // return VD(n, -inf) if the solution doesn't exist // return VD(n, +inf) if the solution is unbounded
11 x,y;
11 w;
                                                                    using VD = vector<double>;
Status mult(const Status& a,const Status& b,ll mod){
                                                                    using VVD = vector<vector<double>>;
                                                                    const double eps = 1e-9;
Status res
 res.x=(a.x*b.x+a.y*b.y%mod*w)%mod;
                                                                    const double inf = 1e+9;
 res.y=(a.x*b.y+a.y*b.x)%mod;
                                                                    int n, m;
 return res;
                                                                    VVD d;
                                                                    vector<int> p, q;
inline Status qpow(Status _base,ll _pow,ll _mod){
                                                                    void pivot(int r, int s) {
 Status res = \{1, 0\};
                                                                      double inv = 1.0 / d[r][s];
 while(_pow>0){
                                                                      for (int i = 0; i < m + 2; ++i)
                                                                      for (int j = 0; j < n + 2; ++j)
if (i != r && j != s)
  if(_pow&1) res=mult(res,_base,_mod);
  _base=mult(_base,_base,_mod);
  _pow>>=1;
                                                                         d[i][j] -= d[r][j] * d[i][s] * inv;
                                                                     for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;
d[r][s] = inv; swap(p[r], q[s]);</pre>
 return res;
inline 11 check(11 x,11 p){
return qpow_mod(x,(p-1)>>1,p);
                                                                    bool phase(int z) {
                                                                      int x = m + z
inline 11 get_root(11 n,11 p){
                                                                      while (true) {
                                                                       int s = -1;
 if(p==2) return 1;
 if(check(n,p)==p-1) return -1;
                                                                       for (int i = 0; i <= n; ++i) {
                                                                        if (!z && q[i] == -1) continue;
 11 a;
 while(true){
                                                                        if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
  a=rand()%p;
  w=((a*a-n)%p+p)%p;
                                                                       if (d[x][s] > -eps) return true;
  if(check(w,p)==p-1) break;
                                                                       int r = -1;
                                                                       for (int i = 0; i < m; ++i) {</pre>
 Status res = \{a, 1\}
                                                                        if (d[i][s] < eps) continue;</pre>
                                                                        if (r == -1 ||
 res=qpow(res,(p+1)>>1,p);
 return res.x;
                                                                         d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
                                                                      if (r == -1) return false;
5.20 De-Bruijn
                                                                      pivot(r, s);
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
 if (t > n) {
                                                                    VD solve(const VVD &a, const VD &b, const VD &c) {
  if (n % p == 0)
                                                                     m = b.size(), n = c.size();
   for (int i = 1; i <= p; ++i)</pre>
                                                                      d = VVD(m + 2, VD(n + 2));
                                                                     for (int i = 0; i < m; ++i)
for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
    res[sz++] = aux[i];
 } else {
  aux[t] = aux[t - p];
                                                                      p.resize(m), q.resize(n + 1);
                                                                      for (int i = 0; i < m; ++i) p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
  db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {
```

for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];

```
q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
for (int i = 1; i < m; ++i)
  if (d[i][n + 1] < d[r][n + 1]) r = i;
if (d[r][n + 1] < -eps) {</pre>
 pivot(r, n);
 if (!phase(1) \mid | d[m + 1][n + 1] < -eps)
   return VD(n, -inf);
 for (int i = 0; i < m; ++i) if (p[i] == -1) {
  int s = min_element(d[i].begin(), d[i].end() - 1)
      - d[i].begin();
  pivot(i, s);
 }
if (!phase(0)) return VD(n, inf);
VD x(n);
for (int i = 0; i < m; ++i)
 if (p[i] < n) x[p[i]] = d[i][n + 1];
return x;
6
    Geometry
6.1
     Basic Geometry
```

```
using coord_t = int;
using Real = double;
using Point = std::complex<coord_t>;
int sgn(coord_t x) {
return (x > 0) - (x < 0);
coord_t dot(Point a, Point b) {
return real(conj(a) * b);
coord_t cross(Point a, Point b) {
return imag(conj(a) * b);
int ori(Point a, Point b, Point c) {
return sgn(cross(b - a, c - a));
bool operator<(const Point &a, const Point &b) {</pre>
return real(a) != real(b)
 ? real(a) < real(b) : imag(a) < imag(b);</pre>
int argCmp(Point a, Point b) {
// -1 / 0 / 1 <-> < / == / > (atan2)
int qa = (imag(a) == 0
   ? (real(a) < 0 ? 3 : 1) : (imag(a) < 0 ? 0 : 2));
int qb = (imag(b) == 0
   ? (real(b) < 0 ? 3 : 1) : (imag(b) < 0 ? 0 : 2));
 if (qa != qb)
 return sgn(qa - qb);
 return sgn(cross(b, a));
template <typename V> Real area(const V & pt) {
coord_t ret = 0;
for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
 ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
return ret / 2.0;
```

#### 6.2 Circle Class

```
struct Circle { Point o; Real r; };
vector<Real> intersectAngle(Circle a, Circle b) {
Real d2 = norm(a.o - b.o);
if (norm(A.r - B.r) >= d2)
 if (A.r < B.r)
  return {-PI, PI};
 else
   return {}:
if (norm(A.r + B.r) <= d2) return {};</pre>
Real dis = hypot(A.x - B.x, A.y - B.y);
Real theta = atan2(B.y - A.y, B.x - A.x);
Real phi = acos((A.r * A.r + d2 - B.r * B.r) /
   (2 * A.r * dis));
Real L = theta - phi, R = theta + phi;
while (L < -PI) L += PI * 2;</pre>
while (R > PI) R -= PI * 2;
return { L, R };
```

```
vector<Point> intersectPoint(Circle a, Circle b) {
 Real d=o.dis(aa.o);
 if (d >= r+aa.r || d <= fabs(r-aa.r)) return {};</pre>
 Real dt = (r*r - aa.r*aa.r)/d, d1 = (d+dt)/2;
 Point dir = (aa.o-o); dir /= d;
 Point pcrs = dir*d1 + o;
 dt=sqrt(max(0.0L, r*r - d1*d1)), dir=dir.rot90();
 return {pcrs + dir*dt, pcrs - dir*dt};
6.3 2D Convex Hull
template<typename PT>
vector<PT> buildConvexHull(vector<PT> d) {
 sort(ALL(d), [](const PT& a, const PT& b){
   return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
 vector<PT> s(SZ(d)<<1);</pre>
 int o = 0;
 for(auto p: d) {
  while(o>=2 && cross(p-s[o-2], s[o-1]-s[o-2])<=0)
   0--:
  s[o++] = p;
 for(int i=SZ(d)-2, t = o+1; i>=0; i--){
  while(o>=t&&cross(d[i]-s[o-2],s[o-1]-s[o-2])<=0)
  s[o++] = d[i];
 s.resize(o-1);
 return s;
6.4 3D Convex Hull
// return the faces with pt indexes
int flag[MXN][MXN];
struct Point{
ld x,y,z;
 Point operator * (const 1d &b) const {
  return (Point) {x*b, y*b, z*b};}
 Point operator * (const Point &b) const {
  return(Point){y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
Point ver(Point a, Point b, Point c) {
return (b - a) * (c - a);}
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0
 REP(i,n) REP(j,n) flag[i][j] = 0;
 vector<Face> now;
 now.emplace_back(0,1,2);
 now.emplace_back(2,1,0)
 for (int i=3; i<n; i++){
  ftop++; vector<Face> next;
  REP(j, SZ(now)) {
   Face& f=now[j]; int ff = 0;
   ld d=(pt[i]-pt[f.a]).dot(
     ver(pt[f.a], pt[f.b], pt[f.c]));
   if (d <= 0) next.push_back(f);</pre>
   if (d > 0) ff=ftop;
   else if (d < 0) ff=-ftop</pre>
   flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
  REP(j, SZ(now)) {
  Face& f=now[j]
   if (flag[f.a][f.b] > 0 &&
     flag[f.a][f.b] != flag[f.b][f.a])
    next.emplace_back(f.a,f.b,i);
   if (flag[f.b][f.c] > 0 &&
     flag[f.b][f.c] != flag[f.c][f.b])
    next.emplace_back(f.b,f.c,i);
   if (flag[f.c][f.a] > 0 &&
     flag[f.c][f.a] != flag[f.a][f.c])
    next.emplace_back(f.c,f.a,i);
 now=next:
 return now;
```

#### 2D Farthest Pair

```
// stk is from convex hull
n = (int)(stk.size());
```

```
for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
int pos = 1, ans = 0; stk.push_back(stk[0]);
                                                                  // Modify p to p_prime
for(int i=0;i<n;i++) {</pre>
 while(abs(cross(stk[i+1]-stk[i],
                                                                  const llf S_prime = calc( p_prime );
                                                                  const llf delta_c = S_prime - S_cur;
llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
   stk[(pos+1)%n]-stk[i]))
   abs(cross(stk[i+1]-stk[i],
   stk[pos]-stk[i]))) pos = (pos+1)%n;
                                                                  if ( rnd( rnd_engine ) <= prob )</pre>
 ans = max({ans, dis(stk[i], stk[pos]),
                                                                   S_cur = S_prime, p = p_prime;
                                                                  if ( S_prime < S_best ) // find min</pre>
  dis(stk[i+1], stk[pos])});
                                                                   S_best = S_prime, p_best = p_prime;
6.6 2D Closest Pair
                                                                 return S_best;
struct cmp_y {
 bool operator()(const P& p, const P& q) const {
                                                                6.9
                                                                     Half Plane Intersection
  return p.y < q.y;</pre>
                                                               // NOTE: Point is complex<Real>
                                                                // cross(pt-line.st, line.dir)<=0 <-> pt in half plane
multiset<P, cmp_y> s;
                                                                struct Line {
void solve(P a[], int n) {
  sort(a, a + n, [](const P& p, const P& q) {
                                                                  Point st, ed;
                                                                  Point dir
  return tie(p.x, p.y) < tie(q.x, q.y);</pre>
                                                                  Line (Point _s, Point _e)
                                                                   : st(_s), ed(_e), dir(_e - _s) {}
 11f d = INF; int pt = 0;
 for (int i = 0; i < n; ++i) {
                                                               bool operator<(const Line &lhs, const Line &rhs) {</pre>
  while (pt < i \text{ and } a[i].x - a[pt].x >= d)
   s.erase(s.find(a[pt++]));
                                                                  if (int cmp = argCmp(lhs.dir, rhs.dir))
  auto it = s.lower_bound(P(a[i].x, a[i].y - d));
                                                                    return cmp == -1;
                                                                  return ori(lhs.st, lhs.ed, rhs.st) < 0;</pre>
  while (it != s.end() and it->y - a[i].y < d)
  d = min(d, dis(*(it++), a[i]));
  s.insert(a[i]);
                                                               Point intersect(const Line &A, const Line &B) {
                                                                  Real t = cross(B.st - A.st, B.dir) /
}
                                                                   cross(A.dir, B.dir);
                                                                  return A.st + t * A.dir;
      kD Closest Pair (3D ver.)
6.7
                                                               }
11f solve(vector<P> v) {
                                                               Real HPI(vector<Line> &lines) {
 shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,</pre>
                                                                  sort(lines.begin(), lines.end());
                                                                  deque<Line> que;
  unordered_map<lld, int>>> m;
                                                                  deque<Point> pt;
 llf d = dis(v[0], v[1]);
 auto Idx = [&d] (llf x) -> lld {
  return round(x * 2 / d) + 0.1; };
                                                                  que.push_back(lines[0]);
                                                                  for (int i = 1; i < (int)lines.size(); i++) {</pre>
                                                                    if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
 auto rebuild_m = [&m, &v, &Idx](int k) {
                                                                     continue
  m.clear();
                                                                #define POP(L, R) \
  for (int i = 0; i < k; ++i)
                                                                    while (pt.size() > 0 \
   m[Idx(v[i].x)][Idx(v[i].y)]
                                                                      && ori(L.st, L.ed, pt.back()) < 0) \
    [Idx(v[i].z)] = i;
                                                                      pt.pop_back(), que.pop_back(); \
 }; rebuild_m(2);
                                                                    while (pt.size() > 0 \
    && ori(R.st, R.ed, pt.front()) < 0) \</pre>
 for (size_t i = 2; i < v.size(); ++i) {</pre>
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
                                                                      pt.pop_front(), que.pop_front();
     kz = Idx(v[i].z); bool found = false;
                                                                    POP(lines[i], lines[i]);
pt.push_back(intersect(que.back(), lines[i]));
  for (int dx = -2; dx <= 2; ++dx) {
   const 11d nx = dx + kx
                                                                    que.push_back(lines[i]);
   if (m.find(nx) == m.end()) continue;
   auto& mm = m[nx];
                                                                  POP(que.front(), que.back())
   for (int dy = -2; dy <= 2; ++dy) {
                                                                  if (que.size() <= 1 ||</pre>
    const 11d ny = dy + ky;
                                                                    argCmp(que.front().dir, que.back().dir) == 0)
    if (mm.find(ny) == mm.end()) continue;
    auto& mmm = mm[ny];
                                                                  pt.push_back(intersect(que.front(), que.back()));
    for (int dz = -2; dz <= 2; ++dz) {
                                                                  return area(pt);
     const 11d nz = dz + kz;
     if (mmm.find(nz) == mmm.end()) continue;
     const int p = mmm[nz];
                                                                6.10 Minkowski sum
     if (dis(v[p], v[i]) < d) {
  d = dis(v[p], v[i]);</pre>
                                                               vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
      found = true;
                                                                 hull(A), hull(B);
                                                                 vector<pll> C(1, A[0] + B[0]), s1, s2;
for(int i = 0; i < SZ(A); ++i)</pre>
     }
                                                                  s1.pb(A[(i + 1) % SZ(A)] - A[i]);
                                                                 for(int i = 0; i < SZ(B); i++)</pre>
                                                                  s2.pb(B[(i + 1) % SZ(B)] - B[i]);
  if (found) rebuild_m(i + 1);
  else m[kx][ky][kz] = i;
                                                                 for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
                                                                  if (p2 >= SZ(B)
                                                                    || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
 return d;
                                                                   C.pb(C.back() + s1[p1++]);
                                                                  else
6.8 Simulated Annealing
                                                                   C.pb(C.back() + s2[p2++]);
                                                                 return hull(C), C;
11f anneal() {
 mt19937 rnd_engine( seed );
 uniform_real_distribution< llf > rnd( 0, 1 );
                                                                6.11 intersection of line and circle
 const llf dT = 0.001;
 // Argument p
                                                               vector<pdd> line_interCircle(const pdd &p1,
 11f S_cur = calc( p ), S_best = S_cur;
                                                                    const pdd &p2,const pdd &c,const double r){
```

```
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 pdd ft=foot(p1,p2,c),vec=p2-p1;
 double dis=abs(c-ft);
 if(fabs(dis-r)<eps) return vector<pdd>{ft};
 if(dis>r) return {};
 vec=vec*sqrt(r*r-dis*dis)/abs(vec);
 return vector<pdd>{ft+vec,ft-vec};
6.12 intersection of polygon and circle
// Divides into multiple triangle, and sum up
// test by HDU2892
const double PI=acos(-1);
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
 if(abs(pb)<eps) return 0;</pre>
 double S, h, theta;
 double a=abs(pb), b=abs(pa), c=abs(pb-pa);
 double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
double cosC = dot(pa,pb) / a / b, C = acos(cosC);
 if(a > r){
 S = (C/2)*r*r;
 h = a*b*sin(C)/c;
  if (h < r && B < PI/2)</pre>
   S = (acos(h/r)*r*r - h*sqrt(r*r-h*h));
else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
    (2 thata)/
  S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
 else S = .5*sin(C)*a*b;
 return S;
double area_poly_circle(const vector<pdd> poly,
 const pdd &0,const double r){
 double S=0;
 for(int i=0;i<SZ(poly);++i)</pre>
  S+=_area(poly[i]-0,poly[(i+1)%SZ(poly)]-0,r)
    *ori(0,poly[i],poly[(i+1)%SZ(poly)]);
 return fabs(S);
6.13 intersection of two circle
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
 pdd o1 = a.0, o2 = b.0;
 double r1 = a.R, r2 = b.R, d2 = abs2(o1 - o2),
     d = sqrt(d2);
 if(d < max(r1, r2) - min(r1, r2) \mid | d > r1 + r2)
  return 0;
 pdd u = (o1 + o2) * 0.5
  + (o1 - o2) * ((r2 * r2 - r1 * r1) / (2 * d2));
 double A = sqrt((r1 + r2 + d) * (r1 - r2 + d)
     * (r1 + r2 - d) * (-r1 + r2 + d));
 pdd v = pdd(o1.Y - o2.Y, -o1.X + o2.X) * A
  / (2 * d2);
 p1 = u + v, p2 = u - v;
 return 1;
6.14 tangent line of two circle
vector<Line> go(const Cir& c1,
  const Cir& c2, int sign1){
 // sign1 = 1 for outer tang, -1 for inter tang
 vector<Line> ret;
 double d_{sq} = norm2(c1.0 - c2.0);
 if( d_sq < eps ) return ret;</pre>
 double d = sqrt( d_sq );
 Pt v = (c2.0 - c1.0) / d;
 double c = (c1.R - sign1 * c2.R) / d;
 if( c * c > 1 ) return ret;
 double h = sqrt( max( 0.0 , 1.0 - c * c ) );
for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
  Pt n = { v.X * c - sign2 * h * v.Y ,
   v.Y * c + sign2 * h * v.X };
  Pt p1 = c1.0 + n * c1.R;
  Pt p2 = c2.0 + n * (c2.R * sign1);
  if( fabs( p1.X - p2.X ) < eps and
    fabs( p1.Y - p2.Y ) < eps )
   p2 = p1 + perp(c2.0 - c1.0);
```

ret.push\_back( { p1 , p2 } );

return ret;

```
6.15
```

```
Minimum Covering Circle
template<typename P>
Circle getCircum(const P &a, const P &b, const P &c){
 Real a1 = a.x-b.x, b1 = a.y-b.y;
 Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
 Real a2 = a.x-c.x, b2 = a.y-c.y;
 Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
 Circle cc;
 cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
 cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2)
 cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
 return cc:
template<typename P>
Circle MinCircleCover(const vector<P>& pts){
 random_shuffle(pts.begin(), pts.end());
 Circle c = { pts[0], 0 };
 for(int i=0;i<(int)pts.size();i++){</pre>
  if (dist(pts[i], c.o) <= c.r) continue;</pre>
  c = \{ pts[i], 0 \};
  for (int j = 0; j < i; j++) {
  if(dist(pts[j], c.o) <= c.r) continue;</pre>
   c.o = (pts[i] + pts[j]) / 2;
   c.r = dist(pts[i], c.o);
for (int k = 0; k < j; k++) {</pre>
    if (dist(pts[k], c.o) <= c.r) continue;</pre>
    c = getCircum(pts[i], pts[j], pts[k]);
 return c;
6.16 KDTree (Nearest Point)
const int MXN = 100005;
struct KDTree {
 struct Node {
  int x,y,x1,y1,x2,y2;
  int id,f;
  Node *L, *R;
 } tree[MXN], *root;
 int n;
 LL dis2(int x1, int y1, int x2, int y2) {
  LL dx = x1-x2, dy = y1-y2;
  return dx*dx+dy*dy;
 static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
 static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 void init(vector<pair<int,int>> ip) {
  n = ip.size();
  for (int i=0; i<n; i++) {</pre>
   tree[i].id = i;
   tree[i].x = ip[i].first;
   tree[i].y = ip[i].second;
  root = build_tree(0, n-1, 0);
 Node* build_tree(int L, int R, int d) {
  if (L>R) return nullptr;
int M = (L+R)/2; tree[M].f = d%2;
  nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
  tree[M].x1 = tree[M].x2 = tree[M].x;
  tree[M].y1 = tree[M].y2 = tree[M].y;
  tree[M].L = build_tree(L, M-1, d+1);
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
   tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
   tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
  tree[M].R = build_tree(M+1, R, d+1);
  if (tree[M].R) {
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
   tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
   tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
  return tree+M;
```

int touch(Node\* r, int x, int y, LL d2){

for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2; for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];

```
LL dis = sqrt(d2)+1;
                                                                    if (uniq) +
  if (x<r->x1-dis || x>r->x2+dis ||
                                                                     for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
    y<r->y1-dis || y>r->y2+dis)
                                                                     return:
   return 0:
                                                                    for (int i = n - 2; i \ge 0; --i)
  return 1;
                                                                     t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
                                                                    pre(sa, c, n, z);
  if (!r || !touch(r, x, y, md2)) return;
                                                                    for (int i = 1; i <= n - 1; ++i)</pre>
  LL d2 = dis2(r->x, r->y, x, y);
                                                                     if (t[i] && !t[i - 1])
  if (d2 < md2 \mid \mid (d2 == md2 \&\& mID < r->id)) {
                                                                       sa[--x[s[i]]] = p[q[i] = nn++] = i;
  mID = r -> id;
                                                                    induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)
   md2 = d2;
                                                                     if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
  }
  // search order depends on split dim
                                                                     bool neq = last < 0 || \</pre>
                                                                      memcmp(s + sa[i], s + last,
(p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
  if ((r->f == 0 \&\& x < r->x) ||
    (r->f == 1 \& y < r->y))
   nearest(r->L, x, y, mID, md2);
nearest(r->R, x, y, mID, md2);
                                                                     ns[q[last = sa[i]]] = nmxz += neq;
  } else {
                                                                    sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
                                                                    pre(sa, c, n, z);
for (int i = nn - 1; i >= 0; --i)
   nearest(r->R, x, y, mID, md2);
   nearest(r\rightarrow L, x, y, mID, md2);
                                                                     sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
                                                                    induce(sa, c, s, t, n, z);
 int query(int x, int y) {
  int id = 1029384756;
                                                                   void build(const string &s) {
  LL d2 = 102938475612345678LL;
                                                                    for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
                                                                     _s[(int)s.size()] = 0; // s shouldn't contain 0
  nearest(root, x, y, id, d2);
                                                                    sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];
for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
  return id;
} tree;
                                                                    int ind = 0; hi[0] = 0;
                                                                    for (int i = 0; i < (int)s.size(); ++i) {
  if (!rev[i]) {</pre>
     Stringology
                                                                      ind = 0;
7.1 Hash
                                                                      continue;
class Hash {
 private:
                                                                     while (i + ind < (int)s.size() && \</pre>
  static constexpr int P = 127, Q = 1051762951;
                                                                       s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  vector<int> h, p;
                                                                     hi[rev[i]] = ind ? ind-- : 0;
  void init(const string &s){
                                                                   }}
   h.assign(s.size()+1, 0); p.resize(s.size()+1);
   for (size_t i = 0; i < s.size(); ++i)
h[i + 1] = add(mul(h[i], P), s[i]);</pre>
                                                                   7.3 Aho-Corasick Algorithm
                                                                   class AhoCorasick{
   generate(p.begin(), p.end(),[x=1,y=1,this]()
                                                                    private:
     mutable{y=x;x=mul(x,P);return y;});
                                                                     static constexpr int Z = 26;
                                                                     struct node{
  int query(int 1, int r){ // 1-base (1, r]
                                                                       node *nxt[ Z ], *fail;
   return sub(h[r], mul(h[1], p[r-1]));}
                                                                       vector< int > data;
                                                                       node(): fail( nullptr ) {
                                                                        memset( nxt, 0, sizeof( nxt ) );
7.2 Suffix Array
                                                                        data.clear();
namespace sfxarray {
                                                                     } *rt;
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
                                                                     inline int Idx( char c ) { return c - 'a'; }
                                                                    public:
int x[maxn], p[maxn], q[maxn * 2];
                                                                     void init() { rt = new node(); }
// sa[i]: sa[i]-th suffix is the \
                                                                     void add( const string& s, int d ) {
// i-th_lexigraphically smallest suffix.
                                                                       node* cur = rt;
// hi[i]: longest common prefix \
                                                                       for ( auto c : s ) {
// of suffix sa[i] and suffix sa[i-1].
                                                                        if ( not cur->nxt[ Idx( c ) ] )
void pre(int *sa, int *c, int n, int z) {
                                                                         cur->nxt[ Idx( c ) ] = new node();
memset(sa, 0, sizeof(int) * n);
                                                                        cur = cur->nxt[ Idx( c ) ];
memcpy(x, c, sizeof(int) * z);
                                                                       cur->data.push_back( d );
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
                                                                     }
memcpy(x + 1, c, sizeof(int) * (z - 1));

for (int i = 0; i < n; ++i)

if (sa[i] && !t[sa[i] - 1])
                                                                     void compile() {
                                                                       vector< node* > bfs;
                                                                       size_t ptr = 0;
                                                                       for ( int i = 0 ; i < Z ; ++ i ) {
   sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
 memcpy(x, c, sizeof(int) * z);
for (int i = n - 1; i >= 0; --i)
                                                                       if ( not rt->nxt[ i ] ) {
                                                                         // uncomment 2 lines to make it DFA
  if (sa[i] && t[sa[i] - 1])
                                                                         // rt->nxt[i] = rt;
   sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
                                                                         continue;
void sais(int *s, int *sa, int *p, int *q,
                                                                        rt->nxt[ i ]->fail = rt;
bool *t, int *c, int n, int z) {
bool uniq = t[n - 1] = true;
                                                                        bfs.push_back( rt->nxt[ i ] );
 int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
                                                                       while ( ptr < bfs.size() ) {</pre>
                                                                       node* u = bfs[ ptr ++ ];
for ( int i = 0 ; i < Z ; ++ i ) {
  if ( not u->nxt[ i ] ) {
 memset(c, 0, sizeof(int) * z);
```

```
// u->nxt[i] = u->fail->nxt[i];
                                                                  break;
      continue:
     node* u_f = u->fail;
                                                               puts(ans ? "Yes" : "No");
     while ( u_f ) {
      if ( not u_f->nxt[ i ] ) {
                                                               return 0;
       u_f = u_f->fail; continue;
                                                              7.5 KMP
      u->nxt[ i ]->fail = u_f->nxt[ i ];
                                                             vector<int> kmp(const string &s) {
      break;
                                                               vector<int> f(s.size(), 0);
     if ( not u_f ) u->nxt[ i ]->fail = rt;
                                                               /* f[i] = length of the longest prefix
     bfs.push_back( u->nxt[ i ] );
                                                                 (excluding s[0:i]) such that it coincides
                                                                 with the suffix of s[0:i] of the same length */
                                                               /* i + 1 - f[i] is the length of the
                                                                 smallest recurring period of s[0:i] */
  void match( const string& s, vector< int >& ret ) {
                                                               int k = 0;
  node* u = rt;
                                                               for (int i = 1; i < (int)s.size(); ++i) {</pre>
                                                               while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
   for ( auto c : s ) {
   while ( u != rt and not u->nxt[ Idx( c ) ] )
                                                                if (s[i] == s[k]) ++k;
    u = u->fail;
                                                               f[i] = k;
    u = u->nxt[Idx(c)];
    if ( not u ) u = rt;
                                                               return f;
    node* tmp = u;
    while ( tmp != rt ) {
                                                             vector<int> search(const string &s, const string &t) {
     for ( auto d : tmp->data )
  ret.push_back( d );
                                                              // return 0-indexed occurrence of t in s
                                                               vector<int> f = kmp(t), r;
     tmp = tmp->fail;
                                                               for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
                                                               while(k > 0 && (k==(int)t.size() \mid \mid s[i]!=t[k]))
                                                                k = f[k - 1]
                                                                if (s[i] == t[k]) ++k;
} ac;
                                                               if (k == (int)t.size()) r.push_back(i-t.size()+1);
     Suffix Automaton
                                                               return res;
struct Node{
Node *green, *edge[26];
                                                              7.6 Z value
int max_len;
                                                             char s[MAXN];
Node(const int _max_len)
  : green(NULL), max_len(_max_len){
                                                             int len,z[MAXN];
  memset(edge, 0, sizeof(edge));
                                                             void Z_value() {
                                                              int i,j,left,right;
} *ROOT, *LAST;
                                                               z[left=right=0]=len;
void Extend(const int c) {
                                                               for(i=1;i<len;i++)</pre>
Node *cursor = LAST;
                                                                j=max(min(z[i-left], right-i),0);
LAST = new Node((LAST->max_len) + 1);
                                                                for(;i+j<len&&s[i+j]==s[j];j++);</pre>
for(;cursor&&!cursor->edge[c]; cursor=cursor->green)
                                                                if(i+(z[i]=j)>right)right=i+z[left=i];
 cursor->edge[c] = LAST;
                                                             }
if (!cursor)
 LAST->green = ROOT;
                                                              7.7
                                                                   Manacher
else {
 Node *potential_green = cursor->edge[c];
                                                             int z[maxn];
                                                              int manacher(const string& s) {
  string t = ".";
  if((potential_green->max_len)==(cursor->max_len+1))
  LAST->green = potential_green;
                                                               for(char c: s) t += c, t += '.';
//assert(potential_green->max_len>(cursor->max_len+1));
                                                               int 1 = 0, r = 0, ans = 0;
                                                               for (int i = 1; i < t.length(); ++i) {
z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
   Node *wish = new Node((cursor->max_len) + 1);
   for(;cursor && cursor->edge[c]==potential_green;
      cursor = cursor->green)
                                                                while (i - z[i] \ge 0 \&\& i + z[i] < t.length()) {
    cursor->edge[c] = wish;
                                                                 if(t[i - z[i]] == t[i + z[i]]) ++z[i];
   for (int i = 0; i < 26; i++)
                                                                 else break;
   wish->edge[i] = potential_green->edge[i];
                                                                if (i + z[i] > r) r = i + z[i], l = i;
   wish->green = potential_green->green;
   potential_green->green = wish;
                                                               for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
   LAST->green = wish;
                                                               return ans;
                                                              7.8 Lexico Smallest Rotation
char S[10000001], A[10000001];
int N;
                                                             string mcp(string s){
                                                              int n = s.length();
int main(){
scanf("%d%s", &N, S);
ROOT = LAST = new Node(0);
                                                               s += s;
                                                              int i=0, j=1;
for (int i = 0; S[i]; i++)
Extend(S[i] - 'a');
                                                              while (i<n && j<n){</pre>
                                                               int k = 0;
                                                                while (k < n \&\& s[i+k] == s[j+k]) k++;
while (N--){
 scanf("%s", A);
                                                               if (s[i+k] <= s[j+k]) j += k+1;</pre>
 Node *cursor = ROOT;
                                                               else i += k+1;
 bool ans = true;
                                                               if (i == j) j++;
  for (int i = 0; A[i]; i++){
   cursor = cursor->edge[A[i] - 'a'];
                                                              int ans = i < n ? i : j;</pre>
   if (!cursor) {
                                                               return s.substr(ans, n);
   ans = false;
```

#### 7.9 BWT

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a
vector<int> v[ SIGMA ];
void BWT(char* ori, char* res){
 // make ori -> ori + ori
  // then build suffix array
void iBWT(char* ori, char* res){
  for( int i = 0 ; i < SIGMA ; i ++ )</pre>
   v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )</pre>
   v[ ori[i] - BASE ].push_back( i );
  vector<int> a:
  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
   for( auto j : v[ i ] ){
  a.push_back( j );
    ori[ ptr ++ ] = BASE + i;
  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
   ptr = a[ ptr ];
  res[ len ] = 0;
} bwt;
```

## 7.10 Palindromic Tree

```
struct palindromic_tree{
struct node{
  int next[26],f,len;
  int cnt, num, st, ed;
 node(int l=0):f(0),len(l),cnt(0),num(0) {
  memset(next, 0, sizeof(next)); }
vector<node> st;
vector<char> s:
int last,n;
void init(){
 st.clear();s.clear();last=1; n=0;
  st.push_back(0);st.push_back(-1);
 st[0].f=1;s.push_back(-1); }
 int getFail(int x){
 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
 return x;}
 void add(int c){
 s.push_back(c-='a'); ++n;
  int cur=getFail(last);
  if(!st[cur].next[c]){
  int now=st.size();
   st.push_back(st[cur].len+2);
  st[now].f=st[getFail(st[cur].f)].next[c];
st[cur].next[c]=now;
   st[now].num=st[st[now].f].num+1;
 last=st[cur].next[c];
 ++st[last].cnt;}
int size(){ return st.size()-2;}
} pt;
int main() {
string s; cin >> s; pt.init();
 for (int i=0; i<SZ(s); i++) {</pre>
 int prvsz = pt.size(); pt.add(s[i]);
 if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [1,r]: s.substr(1, r-1+1)
 }
return 0;
```

#### 8 Misc

### 8.1 Theorems

## 8.1.1 Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii}=d(i)$ ,  $L_{ij}=-c$  where c is the number of edge (i,j) in G.

• The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .

• The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .

### 8.1.2 Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

#### 8.1.3 Cayley's Formula

- Given a degree sequence  $d_1,d_2,\dots,d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .

#### 8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1 + d_2 + \ldots + d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all  $1 \le k \le n$ .

#### 8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

#### 8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let  $N_G(W)$  denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff  $\forall W\subseteq X, |W|\leq |N_G(W)|$ 

## 8.1.7 Euler's planar graph formula

```
V - E + F = C + 1, E \le 3V - 6(?)
```

#### 8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have  $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$ 

## 8.1.9 Lucas's theorem

```
\label{eq:matrix} \binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p}, \text{ where } m=m_k p^k+m_{k-1} p^{k-1}+\cdots+m_1 p+m_0, and n=n_k p^k+n_{k-1} p^{k-1}+\cdots+n_1 p+n_0.
```

## 8.2 MaximumEmptyRect

```
int max_empty_rect(int n, int m, bool blocked[N][N]) {
 static int mxu[2][N], me=0, he=1, ans=0;
 for (int i=0;i<m;i++) mxu[he][i]=0;</pre>
 for (int i=0;i<n;i++) {</pre>
  stack<PII, vector<PII>> stk;
  for (int j=0;j<m;++j) {
  if (blocked[i][j]) mxu[me][j]=0;</pre>
   else mxu[me][j]=mxu[he][j]+1;
   int la = j;
   while (!stk.empty()&&stk.top().FF>mxu[me][j]) {
    int x1 = i - stk.top().FF, x2 = i;
    int y1 = stk.top().SS, y2 = j;
    la = stk.top().SS; stk.pop();
    ans=max(ans, (x2-x1)*(y2-y1));
   if (stk.empty()||stk.top().FF<mxu[me][j])</pre>
    stk.push({mxu[me][j],la});
  while (!stk.empty()) {
   int x1 = i - stk.top().FF, x2 = i;
   int y1 = stk.top().SS-1, y2 = m-1;
   stk.pop(); ans=max(ans,(x2-x1)*(y2-y1));
  swap(me,he);
 return ans;
```

## 8.3 DP-opt Condition

#### 8.3.1 totally monotone (concave/convex)

```
\begin{array}{ll} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

if (it != L.end())

const\_cast<Line\*>(&\*it)->l=now.r=((\*it)&now);

```
8.3.2 monge condition (concave/convex)
                                                                 L.insert(it, now);
\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
                                                                11d Query(11d a) const { // query max at x=a
                                                                 if (L.empty()) return -INF;
8.4 Convex 1D/1D DP
                                                                 Line::flag = false;
struct segment {
                                                                 auto it = --L.upper_bound(\{0, 0, a, 0\});
 int i, 1, r;
                                                                 return (*it)(a);
 segment() {}
 segment(int a, int b, int c): i(a), l(b), r(c) {}
                                                              };
                                                               8.6
                                                                    Josephus Problem
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
                                                               // n people kill m for each turn
 dp[0] = 0;
                                                               int f(int n, int m) {
 deque<segment> dq; dq.push_back(segment(0, 1, n));
                                                                int s = 0;
 for (int i = 1; i <= n; ++i) {
                                                                for (int i = 2; i <= n; i++)
  dp[i] = f(dq.front().i, i);
                                                                 s = (s + m) % i;
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
                                                                return s:
  dq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
                                                               // died at kth
  while (dq.size() &&
                                                               int kth(int n, int m, int k){
   f(i, dq.back().1) < f(dq.back().i, dq.back().1))
                                                                if (m == 1) return n-1;
    dq.pop_back();
                                                                for (k = k+m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  if (dq.size())
                                                                return k;
   int d = 1 << 20, c = dq.back().1;</pre>
   while (d >>= 1) if (c + d <= dq.back().r)</pre>
                                                               8.7 Cactus Matching
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.1 = c + 1;
                                                               vector<int> init_g[maxn],g[maxn*2];
                                                               int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
  if (seg.1 <= n) dq.push_back(seg);</pre>
                                                               void tarjan(int u){
                                                                dfn[u]=low[u]=++dfs_idx;
                                                                for(int i=0;i<(int)init_g[u].size();i++){</pre>
                                                                 int v=init_g[u][i];
     ConvexHull Optimization
8.5
                                                                 if(v==par[u]) continue;
inline 1ld DivCeil(1ld n, 1ld d) { // ceil(n/d)
                                                                 if(!dfn[v]){
 return n / d + (((n < 0) != (d > 0)) \&\& (n % d));
                                                                  par[v]=u
                                                                  tarjan(v);
struct Line {
                                                                  low[u]=min(low[u],low[v]);
 static bool flag;
                                                                  if(dfn[u]<low[v]){</pre>
11d a, b, 1, r; // y=ax+b in [1, r)
11d operator()(11d x) const { return a * x + b; }
                                                                   g[u].push_back(v);
                                                                   g[v].push_back(u);
 bool operator<(const Line& i) const {</pre>
                                                                 }else{
  return flag ? tie(a, b) < tie(i.a, i.b) : 1 < i.l;</pre>
                                                                  low[u]=min(low[u],dfn[v]);
 11d operator&(const Line& i) const {
                                                                  if(dfn[v]<dfn[u]){</pre>
                                                                   int temp_v=u;
  return DivCeil(b - i.b, i.a - a);
                                                                   bcc_id++;
                                                                   while(temp_v!=v){
bool Line::flag = true;
                                                                    g[bcc_id+n].push_back(temp_v);
class ConvexHullMax {
                                                                    g[temp_v].push_back(bcc_id+n);
 set<Line> L;
                                                                    temp_v=par[temp_v];
 public:
 ConvexHullMax() { Line::flag = true; }
                                                                   g[bcc_id+n].push_back(v);
 void InsertLine(lld a, lld \bar{b}) { // add y = ax + b
                                                                   g[v].push_back(bcc_id+n);
  Line now = \{a, b, -INF, INF\};
                                                                   reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
  if (L.empty()) {
   L.insert(now);
   return;
  Line::flag = true;
                                                               int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
  auto it = L.lower_bound(now);
                                                               void dfs(int u,int fa){
                                                                if(u<=n){</pre>
  auto prv = it == L.begin() ? it : prev(it);
  if (it != L.end() && ((it != L.begin() &&
                                                                 for(int i=0;i<(int)g[u].size();i++){</pre>
   (*it)(it->1) >= now(it->1) &&
                                                                  int v=g[u][i];
   (*prv)(prv->r - 1) >= now(prv->r - 1)) ||
                                                                  if(v==fa) continue;
   (it == L.begin() && it->a == now.a))) return;
                                                                  dfs(v,u);
  if (it != L.begin()) {
                                                                  memset(tp,0x8f,sizeof tp);
   while (prv != L.begin() &&
                                                                  if(v<=n){
    (*prv)(prv->1) <= now(prv->1))
                                                                   tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
                                                                   tp[1]=max(
     prv = --L.erase(prv)
   if (prv == L.begin() && now.a == prv->a)
                                                                    dp[u][0]+dp[v][0]+1
    L.erase(prv);
                                                                    dp[u][1]+max(dp[v][0],dp[v][1])
  if (it != L.end())
                                                                  }else{
   while (it != --L.end() &&
                                                                   tp[0]=dp[u][0]+dp[v][0];
    (*it)(it->r) <= now(it->r))
                                                                   tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
     it = L.erase(it);
  if (it != L.begin()) {
                                                                  dp[u][0]=tp[0],dp[u][1]=tp[1];
   prv = prev(it);
                                                                 }
   const_cast<Line*>(&*prv)->r=now.l=((*prv)&now);
                                                                }else{
                                                                 for(int i=0;i<(int)g[u].size();i++){</pre>
```

int v=g[u][i];

if(v==fa) continue;

```
dfs(v,u);
                                                                  for(int j = L[i]; j != i; j = L[j]) {
                                                                   U[D[j]] = j;
 min_dp[0][0]=0;
                                                                   D[U[j]] = j
  min_dp[1][1]=1;
                                                                   ++S[col[j]];
  min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
  for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                void dance(int d) {
   int v=g[u][i];
   if(v==fa) continue;
                                                                 if(d>=ansd) return;
   memset(tmp,0x8f,sizeof tmp);
                                                                 if(R[0] == 0) {
   tmp[0][0]=max(
                                                                  ansd = d;
   min_dp[0][0]+max(dp[v][0],dp[v][1]),
                                                                  return:
    min_dp[0][1]+dp[v][0]
                                                                 int c = R[0];
                                                                 for(int i = R[0]; i; i = R[i])
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
   tmp[1][0]=max(
                                                                  if(S[i] < S[c]) c = i;
   \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
                                                                 remove(c);
                                                                 for(int i = D[c]; i != c; i = D[i]) {
    min_dp[1][1]+dp[v][0]
                                                                  ans[d] = row[i]
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
                                                                  for(int j = R[i]; j != i; j = R[j])
  memcpy(min_dp,tmp,sizeof tmp);
                                                                   remove(col[j]);
                                                                  dance(d+1);
  dp[u][1]=max(min_dp[0][1], min_dp[1][0]);
                                                                  for(int j = L[i]; j != i; j = L[j])
  dp[u][0]=min_dp[0][0];
                                                                   resume(col[j]);
                                                                 resume(c);
int main(){
                                                                }
                                                               } sol;
int m,a,b;
scanf("%d%d",&n,&m);
for(int i=0;i<m;i++){
  scanf("%d%d",&a,&b);</pre>
                                                               8.9 Tree Knapsack
                                                               int dp[N][K];PII obj[N];
  init_g[a].push_back(b);
                                                               vector<int> G[N];
 init_g[b].push_back(a);
                                                               void dfs(int u, int mx){
                                                                for(int s: G[u]) {
par[1]=-1;
                                                                 if(mx < obj[s].first) continue;</pre>
tarjan(1);
                                                                 for(int i=0;i<=mx-obj[s].FF;i++)</pre>
dfs(1,-1);
                                                                  dp[s][i] = dp[u][i];
printf("%d\n", max(dp[1][0], dp[1][1]));
                                                                 dfs(s, mx - obj[s].first);
return 0;
                                                                 for(int i=obj[s].FF;i<=mx;i++)</pre>
                                                                  dp[u][i] = max(dp[u][i],
                                                                   dp[s][i - obj[s].FF] + obj[s].SS);
8.8 DLX
                                                                }
struct DLX {
const static int maxn=210;
                                                               int main(){
const static int maxm=210;
                                                                int n, k; cin >> n >> k;
const static int maxnode=210*210;
                                                                for(int i=1;i<=n;i++){</pre>
 int n, m, size, row[maxnode], col[maxnode];
                                                                 int p; cin >> p;
int U[maxnode], D[maxnode], L[maxnode], R[maxnode];
                                                                 G[p].push_back(i);
int H[maxn], S[maxm], ansd, ans[maxn];
                                                                 cin >> obj[i].FF >> obj[i].SS;
void init(int _n, int _m) {
 n = _n, m = _m;
                                                                dfs(0, k); int ans = 0;
  for(int i = 0; i <= m; ++i) {</pre>
                                                                for(int i=0; i<=k; i++) ans = max(ans, dp[0][i]);
   S[i] = 0;
                                                                cout << ans << '\n';
  U[i] = D[i] = i;
                                                                return 0;
  L[i] = i-1, R[i] = i+1;
 R[L[0] = size = m] = 0;
                                                               8.10 N Queens Problem
  for(int i = 1; i <= n; ++i) H[i] = -1;
                                                               vector< int > solve( int n ) {
                                                                // no solution when n=2, 3
 void Link(int r, int c) {
 ++S[col[++size] = c];
                                                                vector< int > ret;
                                                                if ( n % 6 == 2 ) {
for ( int i = 2 ; i <= n ; i += 2 )
 row[size] = r; D[size] = D[c];
U[D[c]] = size; U[size] = c; D[c] = size;
  if(H[r] < 0) H[r] = L[size] = R[size] = size;</pre>
                                                                  ret.push_back( i );
                                                                 ret.push_back( 3 ); ret.push_back( 1 );
for ( int i = 7 ; i <= n ; i += 2 )
  else {
  R[size] = R[H[r]];
   L[R[H[r]]] = size;
                                                                  ret.push_back( i );
                                                                ret.push_back( 5 );
} else if ( n % 6 == 3 ) {
   L[size] = H[r];
   R[H[r]] = size;
                                                                 for ( int i = 4 ; i <= n ; i += 2 )
 }
                                                                  ret.push_back( i );
void remove(int c) {
  L[R[c]] = L[c]; R[L[c]] = R[c];
                                                                 ret.push_back( 2 );
                                                                 for ( int i = 5 ; i <= n ; i += 2 )
                                                                  ret.push_back( i );
  for(int i = D[c]; i != c; i = D[i])
                                                                 ret.push_back( 1 ); ret.push_back( 3 );
   for(int j = R[i]; j != i; j = R[j]) {
   U[D[j]] = U[j];
                                                                } else {
                                                                 for ( int i = 2 ; i <= n ; i += 2 )
    D[U[j]] = D[j];
                                                                  ret.push_back( i );
    --S[col[j]];
                                                                 for ( int i = 1 ; i <= n ; i += 2 )
                                                                  ret.push_back( i );
void resume(int_c) {
 L[R[c]] = c; R[L[c]] = c;
                                                                return ret;
 for(int i = U[c]; i != c; i = U[i])
                                                               }
```

## 8.11 Aliens Optimization

```
long long Alien() {
long long c = kInf;
for (int d = 60; d >= 0; --d) {
   // cost can be negative, depending on the problem.
   if (c - (1LL << d) < 0) continue;
   long long ck = c - (1LL << d);
   pair<long long, int> r = check(ck);
   if (r.second == k) return r.first - ck * k;
   if (r.second < k) c = ck;
}
pair<long long, int> r = check(c);
return r.first - c * k;
}
```