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4.1 Kuhn Munkres 4.2 Bipartize Matching 5.3 General Graph Matching 6.3 General Graph Matching 6.4 Hinimum Weight Matching (Clique version) 6.5 Flow Models 6.5 Flow Models 6.6 Climate Tollar State Climate State State Climate State Climate State Climate State Climate State State Climate State Climate State Climate State Climate State Climate State State State State State State State State State Climat	3.11Minimum Steiner Tree	<pre>8 syn on 8 colorscheme ron 9 filetype indent on map <f8> <esc>:w<cr>:!clear && g++ "%" -o "%<" -</cr></esc></f8></pre>
4.8 Global Min-Cut 5 Math 5.1 Prime Table 5.2 [*] Enumeration 5.3 axthy-god 5.5 P.1 Count (Linear Sieve) 5.5 P.1 Count (Linear Sieve) 5.6 Range Sieve 5.7 Miller Rabin 5.8 Inverse Element 5.1 Speared Phi Function 5.1 Speared Phi Function 5.1 Speared Phi Function 5.1 Speared Reasinder 5.2 Speared Reasinder 5.3 Speared Reasinder 5.4 Speared Reasinder 5.5 Triangle Coccurrent Reasing	4.1 Kuhn Munkres	success <cr> map <f9> <esc>:w<cr>:!clear && g++ "%" -o "%<" -02 && echo success<cr> map <f10> <esc>:!./"%<"<cr></cr></esc></f10></cr></cr></esc></f9></cr>
5.1 Prime Table 12	4.7 Minimum Cost Maximum Flow	11
1.3 1.3	5.1 Prime Table	#pragma GCC optimize("orast, no-stack-protector) #pragma GCC optimize("no-math-errno, unroll-loops") #pragma GCC target("sse, sse2, sse3, ssse3, sse4") #pragma GCC target("popcnt, abm, mmx, avx, tune=native")
S. 10Gauss Elimination	5.6 Range Sieve	13 1.3 IO Optimization 13 static inline int gc() {
S. 16Polynomial Operations	5.10Gauss Elimination	<pre>static char buf[1 << 20], *p = buf, *end = buf; if (p == end) { end = buf + fread(buf, 1, 1 << 20, stdin); if (end == buf) return EOF;</pre>
S-21Simplex Construction	5.16Polynomial Operations	
17 6.1 Point Class 17 6.2 Circle Class 18 6.3 Segment Class 18 6.4 Line Class 18 6.5 Triangle Circumcentre 19 6.6 2D Convex Hull 19 6.7 2D Farthest Pair 19 6.8 2D Closest Pair 19 6.9 Simulated Annealing 19 6.10Half Plane Intersection 19 6.11Ternary Search on Integer 6.12Minimum Covering Circle 20 6.13KDTree (Nearest Point) 20 7.2 Suffix Array 21 7.3 Aho-Corasick Algorithm 21 7.4 Suffix Automaton 21 7.5 KMP 22 7.6 Z value 7.7 Manacher 22 7.8 Lexico Smallest Rotation 22 7.9 BWT 22 Static Constexpr int LOG_BASE = 9;	5.21Simplex Construction	register int c = gc(); register T = 1; _ = 0; while(('0'>c c>'9') && c!=EOF && c!='-') c = gc();
6.6 2D Convex Hull 6.7 2D Farthest Pair 6.8 2D Closest Pair 6.9 Simulated Annealing 6.9 Simulated Annealing 6.10 Half Plane Intersection 6.10 Half Plane Intersection 6.12 Minimum Covering Circle 6.13 KDTree (Nearest Point) 7 Stringology 7.1 Hash 7.2 Suffix Array 7.3 Aho-Corasick Algorithm 7.4 Suffix Automaton 7.5 KMP 7.6 Z value 7.7 Manacher 7.8 Lexico Smallest Rotation 7.9 BWT 7.1 BASE = 100000000000000000000000000000000000	6.1 Point Class	<pre>if(c == EOF) return false; while('0'<=c&&c<='9') _ = _ * 10 + c - '0', c = gc();</pre>
6.11Ternary Search on Integer 6.12Minimum Covering Circle 6.13KDTree (Nearest Point) 7 Stringology 7.1 Hash 7.2 Suffix Array 7.3 Aho-Corasick Algorithm 7.4 Suffix Automaton 7.5 KMP 7.6 Z value 7.7 Manacher 7.7 Manacher 7.8 Lexico Smallest Rotation 7.9 BWT 7.9 Lexico Smallest Rotation 20 20 2 Data Structure 2.1 Class BigInt 2.1 private: using 1ld = int_fast64_t; #define PRINTF_ARG PRIdFAST64 #define LOG_BASE_STR "9" static constexpr 1ld BASE = 1000000000; static constexpr int LOG_BASE = 9;	6.6 2D Convex Hull	template < typename T, typenameArgs > static inline bool gn(T &x, Args &args) { return gn(x) && gn(args); }
7.1 Hash	6.11Ternary Search on Integer	20 2 Data Structure
7.3 Aho-Corasick Algorithm	7.1 Hash	20
7.0 2 Value	7.3 Aho-Corasick Algorithm	private: using lld = int_fast64_t; #define PRINTE ARG PRIJEAST64
A TREATMORPHUR TRAD	7.7 Manacher	#define LOG_BASE_STR "9" static constexpr lld BASE = 1000000000;

BigInt ret; ret.dig.resize(len()+a.len()+1);

```
inline int len() const { return (int) dig.size(); }
                                                                 ret.neg = neg ^ a.neg;
                                                                 for(int i=0;i<len();i++)</pre>
 inline int cmp_minus(const BigInt& a) const {
   if(len() == 0 && a.len() == 0) return 0;
                                                                   for(int j=0;j<a.len();j++){</pre>
                                                                     ret.dig[i+j] += dig[i] * a.dig[j];
   if(neg ^ a.neg)return a.neg ^ 1;
   if(len()!=a.len())
                                                                     if(ret.dig[i+j] >= BASE) {
                                                                       11d x = ret.dig[i+j] / BASE;
     return neg?a.len()-len():len()-a.len();
   for(int i=len()-1;i>=0;i--) if(dig[i]!=a.dig[i])
                                                                       ret.dig[i+j+1] += x
                                                                       ret.dig[i+j] -= x * BASE;
     return neg?a.dig[i]-dig[i]:dig[i]-a.dig[i];
   return 0:
 inline void trim(){
                                                                 ret.trim(); return ret;
   while(!dig.empty()&&!dig.back())dig.pop_back();
   if(dig.empty()) neg = false;
                                                               BigInt operator/(const BigInt& a) const {
                                                                 assert(a.len());
public:
                                                                 if(len() < a.len()) return 0;</pre>
 BigInt(): dig(vector<lld>()), neg(false){}
                                                                 BigInt ret; ret.dig.resize(len()-a.len()+1);
 BigInt(lld a): dig(vector<lld>())
                                                                 ret.neg = a.neg;
                                                                 for(int i=len()-a.len();i>=0;i--){
   neg = a<0; dig.push_back(abs(a));</pre>
   trim();
                                                                   11d 1 = 0, r = BASE;
                                                                   while (r-1 > 1) {
                                                                     11d \ mid = (1+r)>>1;
 BigInt(const string& a): dig(vector<lld>()){
   assert(!a.empty()); neg = (a[0]=='-');
                                                                     ret.dig[i] = mid;
                                                                     if(ret*a<=(neg?-(*this):(*this))) 1 = mid;</pre>
   for(int i=((int)a.size())-1;i>=neg;i-=LOG_BASE){
     11d cur = 0;
                                                                     else r = mid;
     for(int j=min(LOG_BASE-1,i-neg);j>=0;j--)
       cur = cur*10+a[i-j]-'0';
                                                                   ret.dig[i] = 1;
     dig.push_back(cur);
                                                                 ret.neg ^= neg; ret.trim();
   } trim();
                                                                 return ret;
 inline bool operator<(const BigInt& a)const
                                                               BigInt operator%(const BigInt& a) const {
   {return cmp_minus(a)<0;}
 inline bool operator<=(const BigInt& a)const
                                                                 return (*this) - (*this) / a * a;
   {return cmp_minus(a)<=0;}
                                                               friend BigInt abs(BigInt a) { a.neg = 0; return a; }
 inline bool operator==(const BigInt& a)const
   {return cmp_minus(a)==0;}
                                                               friend void swap(BigInt& a, BigInt& b){
 inline bool operator!=(const BigInt& a)const
                                                                 swap(a.dig, b.dig); swap(a.neg, b.neg);
   {return cmp_minus(a)!=0;}
 inline bool operator>(const BigInt& a)const
                                                               friend istream& operator>>(istream& ss, BigInt& a){
                                                                 string s; ss >> s; a = s; return ss;
   {return cmp_minus(a)>0;}
 inline bool operator>=(const BigInt& a)const
   {return cmp_minus(a)>=0;}
                                                               friend ostream&operator<<(ostream&o, const BigInt&a){</pre>
                                                                 if(a.len() == 0) return o << '0';
if(a.neg) o << '-';
 BigInt operator-() const {
   BigInt ret = *this;
   ret.neg ^= 1; return ret;
                                                                 ss << o.dig.back();
                                                                 for(int i=a.len()-2;i>=0;i--)
 BigInt operator+(const BigInt& a) const {
                                                                   o<<setw(LOG_BASE)<<setfill('0')<<a.dig[i];
   if(neg) return -(-(*this)+(-a));
                                                                 return o;
   if(a.neg) return (*this)-(-a);
   int n = max(a.len(), len());
                                                               inline void print() const {
                                                                 if(len() == 0){putchar('0');return;}
if(neg) putchar('-');
printf("%" PRINTF_ARG, dig.back());
   BigInt ret; ret.dig.resize(n);
   11d pro = 0;
   for(int i=0;i<n;i++) {</pre>
     ret.dig[i] = pro;
                                                                 for(int i=len()-2;i>=0;i--)
     if(i < a.len()) ret.dig[i] += a.dig[i];</pre>
                                                                   printf("%0" LOG_BASE_STR PRINTF_ARG, dig[i]);
     if(i < len()) ret.dig[i] += dig[i];</pre>
     pro = 0:
                                                               #undef PRINTF ARG
     if(ret.dig[i] >= BASE) pro = ret.dig[i]/BASE;
                                                               #undef LOG_BASE_STR
     ret.dig[i] -= BASE*pro;
   if(pro != 0) ret.dig.push_back(pro);
                                                            2.2 Dark Magic
   return ret;
 BigInt operator-(const BigInt& a) const {
                                                            #include <ext/pb_ds/assoc_container.hpp>
   if(neg) return -(-(*this) - (-a));
                                                             #include <ext/pb_ds/priority_queue.hpp>
   if(a.neg) return (*this) + (-a);
                                                             using __gnu_pbds::pairing_heap_tag;
   int diff = cmp_minus(a);
                                                             using __gnu_pbds::binary_heap_tag;
   if(diff < 0) return -(a - (*this));</pre>
                                                             using __gnu_pbds::binomial_heap_tag;
                                                             using __gnu_pbds::rc_binomial_heap_tag;
   if(diff == 0) return 0;
   BigInt ret; ret.dig.resize(len(), 0);
                                                             using __gnu_pbds::thin_heap_tag;
   for(int i=0;i<len();i++) {</pre>
                                                             template<typename T>
     ret.dig[i] += dig[i];
                                                             using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>,\
     if(i < a.len()) ret.dig[i] -= a.dig[i];</pre>
                                                                                                    pairing_heap_tag>;
     if(ret.dig[i] < 0){</pre>
                                                             // a.join(b), pq.modify(pq.push(10), 87)
       ret.dig[i] += BASE;
                                                             using __gnu_pbds::rb_tree_tag;
                                                             using __gnu_pbds::ov_tree_tag;
       ret.dig[i+1]--;
                                                             using __gnu_pbds::splay_tree_tag;
                                                             template<typename T>
                                                            using ordered_set = __gnu_pbds::tree<T,\
__gnu_pbds::null_type,less<T>,rb_tree_tag,\
   ret.trim(); return ret;
 BigInt operator*(const BigInt& a) const {
                                                             __gnu_pbds::tree_order_statistics_node_update>;
   if(!len()||!a.len()) return 0;
                                                             // find_by_order, order_of_key
```

template<typename A, typename B>

inline void rotate(Node* node){

Node* par=node->par;

```
using hTable1=__gnu_pbds::cc_hash_table<A,B>;
                                                               Node* par_par=par->par
                                                               bool dir=node->is_rch();
template<typename A, typename B>
using hTable2=__gnu_pbds::gp_hash_table<A,B>;
                                                               bool par_dir=par->is_rch();
                                                               to_child(par, node->ch[!dir], dir);
                                                               to_child(node,par,!dir);
2.3 Disjoint Set
                                                               if(par_par!=nullptr && par_par->ch[par_dir]==par)
                                                                 to_child(par_par,node,par_dir);
class DJS {
                                                               else node->par=par_par;
private:
  vector< int > fa, sz, sv;
                                                             inline void splay(Node* node){
  vector< pair< int*, int > > opt;
                                                               Node* tmp=node:
  void assign( int *k, int v ) {
                                                               stk[top++]=node;
    opt.emplace_back( k, *k );
                                                               while(!tmp->is_root()){
    *k = v;
                                                                 tmp=tmp->par;
                                                                 stk[top++]=tmp;
public:
  void init( int n ) {
                                                               while(top) stk[--top]->down();
for(Node *fa=node->par;
    fa.resize( n ); iota( fa.begin(), fa.end(), 0 );
sz.resize( n ); fill( sz.begin(), sz.end(), 1 );
                                                                !node->is_root();
    opt.clear();
                                                                rotate(node), fa=node->par)
                                                                 if(!fa->is_root())
  int query(int x) {return fa[x] == x?x:query(fa[x]);}
                                                                   rotate(fa->is_rch()==node->is_rch()?fa:node);
  void merge( int a, int b ) {
    int af = query( a ), bf = query( b );
                                                             inline void access(Node* node){
    if( af == bf ) return;
                                                               Node* last=nullptr;
    if( sz[ af ] < sz[ bf ] ) swap( af, bf );</pre>
                                                               while(node!=nullptr){
    assign( &fa[ bf ], fa[ af ] );
assign( &sz[ af ], sz[ af ] + sz[ bf ] );
                                                                 splay(node);
                                                                 to_child(node, last, true);
                                                                 last=node;
  void save() { sv.push_back( (int) opt.size() ); }
                                                                 node=node->par;
  void undo() {
    int ls = sv.back(); sv.pop_back();
    while ( ( int ) opt.size() > ls )
                                                             inline void change_root(Node* node){
      pair< int*, int > cur = opt.back();
                                                               access(node);splay(node);node->set_rev();
      .
*cur.first = cur.second;
      opt.pop_back();
                                                             inline void link(Node* x, Node* y){
                                                               change_root(x);splay(x);x->par=y;
  }
};
                                                             inline void split(Node* x, Node* y){
                                                               change_root(x);access(y);splay(x);
      Link-Cut Tree
2.4
                                                               to_child(x,nullptr,true);y->par=nullptr;
                                                             inline void change_val(Node* node,int v){
struct Node{
                                                               access(node);splay(node);node->v=v;node->up();
  Node *par,*ch[2];
  int xor_sum,v;
  bool is_rev;
                                                             inline int query(Node* x, Node* y){
                                                               change_root(x);access(y);splay(y);
  Node(int _v){
    v=xor_sum=_v;is_rev=false;
                                                               return y->xor_sum;
    par=ch[0]=ch[1]=nullptr;
                                                             inline Node* find_root(Node* node){
  inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
                                                               access(node);splay(node);
  inline void down(){
                                                               Node* last=nullptr
    if(is_rev)
                                                               while(node!=nullptr){
                                                                 node->down();last=node;node=node->ch[0];
      if(ch[0]!=nullptr) ch[0]->set_rev();
      if(ch[1]!=nullptr) ch[1]->set_rev();
                                                               return last;
      is rev=false:
    }
                                                             set<pii> dic;
  inline void up(){
                                                             inline void add_edge(int u,int v){
                                                               if(u>v) swap(u,v);
    xor_sum=v;
    if(ch[0]!=nullptr){
                                                               if(find_root(node[u])==find_root(node[v])) return;
                                                               dic.insert(pii(u,v))
      xor_sum^=ch[0]->xor_sum;
      ch[0]->par=this;
                                                               link(node[u],node[v]);
    if(ch[1]!=nullptr){
                                                             inline void del_edge(int u,int v){
      xor_sum^=ch[1]->xor_sum;
                                                               if(u>v) swap(u,v);
      ch[1]->par=this;
                                                               if(dic.find(pii(u,v))==dic.end()) return;
                                                               dic.erase(pii(u,v))
                                                               split(node[u], node[v]);
  inline bool is_root(){
    return par==nullptr ||\
      (par->ch[0]!=this && par->ch[1]!=this);
                                                             2.5 LiChao Segment Tree
  bool is_rch(){return !is_root() && par->ch[1]==this;}
struct Line{
                                                               int m, k, id;
void to_child(Node* p,Node* c,bool dir){
                                                               Line() : id( -1 ) {}
                                                               Line( int a, int b, int c )
: m( a ), k( b ), id( c ) {}
  p->ch[dir]=c;
  p->up();
```

int at(int x) { return m * x + k; }

class LiChao {

```
private:
    int n; vector< Line > nodes;
    inline int lc( int x ) { return 2 * x + 1; }
    inline int rc( int x ) { return 2 * x + 2;
    void insert( int 1, int r, int id, Line ln ) {
      int m = (1 + r) >> 1;
      if ( nodes[ id ].id == -1 ) {
        nodes[ id ] = ln;
        return:
      bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
      if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
        atLeft ^= 1; swap( nodes[ id ], ln );
      if ( r - 1 == 1 ) return;
      if ( atLeft ) insert( 1, m, lc( id ), ln );
      else insert( m, r, rc( id ), ln );
    int query( int 1, int r, int id, int x ) {
      int ret = 0;
      if ( nodes[ id ].id != -1 )
        ret = nodes[ id ].at( x );
      int m = ( 1 + r ) >> 1;
if ( r - 1 == 1 ) return ret;
      else if (x < m)
        return max( ret, query( 1, m, lc( id ), x ) );
      else
        return max( ret, query( m, r, rc( id ), x ) );
  public:
    void build( int n_ ) {
      n = n_; nodes.clear();
      nodes.resize( n << 2, Line() );</pre>
    void insert( Line ln ) { insert( 0, n, 0, ln ); }
    int query( int x ) { return query( 0, n, 0, x ); }
} lichao;
```

2.6 Treap

```
namespace Treap{
 #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
 struct node{
    int size;
    uint32_t pri;
   node *lc, *rc;
    node() : size(0), pri(rand()), lc( 0 ), rc( 0 ) {}
    void pull() {
     size = 1
     if ( lc ) size += lc->size;
     if ( rc ) size += rc->size;
   }
 };
 node* merge( node* L, node* R ) {
    if ( not L or not R ) return L ? L : R;
    if ( L->pri > R->pri ) {
     L->rc = merge( L->rc, R ); L->pull();
     return L;
    } else {
     R->lc = merge( L, R->lc ); R->pull();
      return R;
   }
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
   if ( not rt ) L = R = nullptr;
    else if( sz( rt->lc ) + 1 <= k ) {
     L = rt;
      split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
     L->pull();
    } else {
     R = rt;
      split_by_size( rt->lc, k, L, R->lc );
     R->pull();
    }
  #undef sz
```

2.7 SparseTable

```
template < typename T, typename Cmp_ = less< T > >
class SparseTable {
private:
  vector< vector< T > > tbl;
  vector< int > lg;
  T cv(Ta, Tb) {
    return Cmp_()( a, b ) ? a : b;
public:
  void init( T arr[], int n ) {
     // 0-base
    lg.resize(n+1);
    lg[0] = -1;
     for( int i=1 ; i<=n ; ++i ) lg[i] = lg[i>>1] + 1;
    tbl.resize(lg[n] + 1);
    tbl[ 0 ].resize( n );
    copy( arr, arr + n, tbl[ 0 ].begin() );
for ( int i = 1 ; i <= lg[ n ] ; ++ i ) {</pre>
      int len = 1 << ( i - 1 ), sz = 1 << i;
       tbl[ i ].resize( n - sz + 1 );
      for ( int j = 0 ; j <= n - sz ;
         tbl[i][j] = cv(tbl[i-1][j], tbl[i-1][j+len]);
  T query( int l, int r ) {
    // 0-base [1, r)
    int wh = \lg[r-1],
                            len = 1 << wh;
    return cv( tbl[ wh ][ 1 ], tbl[ wh ][ r - len ] );
};
```

2.8 Linear Basis

```
struct LinearBasis {
private:
  int n, sz
  vector< llu > B;
  inline llu two( int x ){ return ( ( llu ) 1 ) << x; }</pre>
public:
  void init( int n_ ) {
    n = n_{;} B.clear(); B.resize(n); sz = 0;
  void insert( llu x ) {
     // add x into B
    for ( int i = n-1; i >= 0 ; --i ) if( two(i) & x ){
  if ( B[ i ] ) x ^= B[ i ];
       else {
         B[ i ] = x; sz++;
for ( int j = i - 1 ; j >= 0 ; -- j
           if( B[ j ] && ( two( j ) & B[ i ] ) )
B[ i ] ^= B[ j ];
         for (int j = i + 1)
           if ( two( i ) & B[ j ] )
  B[ j ] ^= B[ i ];
         break:
      }
    }
  inline int size() { return sz; }
  bool check( llu x ) {
     // is x in span(B) ?
    for ( int i = n-1 ; i >= 0 ; --i ) if( two(i) & x )
       if( B[ i ] ) x ^= B[ i ];
       else return false;
    return true;
  llu kth_small(llu k) {
     /** 1-base would always > 0 **/
     /** should check it **/
     /* if we choose at least one element
        but size(B)(vectors in B)==N(original elements)
        then we can't get 0 */
    llu ret = 0;
     for ( int i = 0 ; i < n ; ++ i ) if( B[ i ] ) {
      if( k & 1 ) ret ^= B[ i ];
       k >>= 1;
    return ret:
} base;
```

3 Graph

3.1 Euler Circuit

```
bool vis[ N ]; size_t la[ K ];
void dfs( int u, vector< int >& vec ) {
  while ( la[ u ] < G[ u ].size() ) {
    if( vis[ G[ u ][ la[ u ] ].second ] ) {
        ++ la[ u ];
        continue;
    }
    int v = G[ u ][ la[ u ] ].first;
    vis[ G[ u ][ la[ u ] ].second ] = true;
    ++ la[ u ]; dfs( v, vec );
    vec.push_back( v );
}</pre>
```

3.2 BCC Edge

```
class BCC{
private:
  vector< int > low, dfn;
  int cnt;
  vector< bool > bridge;
  vector< vector< PII > > G;
  void dfs( int w, int f ) {
    low[ w ] = dfn[ w ] = cnt++;
    for ( auto [ u, t ] : G[ w ] ) {
      if ( u == f ) continue;
if ( dfn[ u ] != 0 ) {
        low[ w ] = min( low[ w ], dfn[ u ] );
      }else{
        dfs( u, w );
        low[ w ] = min( low[ w ], low[ u ] );
        if ( low[ u ] > dfn[ w ] ) bridge[ t ] = true;
    }
public:
  void init( int n, int m ) {
    G.resize(n); cnt = 0;
    fill( G.begin(), G.end(), vector< PII >() );
    bridge.clear(); bridge.resize( m );
    low.clear(); low.resize( n );
dfn.clear(); dfn.resize( n );
  void add_edge( int u, int v ) {
    // should check for multiple edge
    G[ u ].emplace_back( v, cnt );
    G[ v ].emplace_back( u, cnt ++ );
  void solve(){
    cnt = 1;
    for (int i = 0; i < n; ++i)</pre>
      if (not vis[ i ]) dfs(i, i);
  // the id will be same as insert order, \theta-base
  bool is_bridge( int x ) { return bridge[ x ]; }
} bcc;
```

3.3 BCC Vertex

```
class BCC {
 private:
    int n, t, ecnt;
    vector<vector<pair<int, int>>> G;
    vector<int> low, tin, st, bcc;
    vector<bool> ap, ins;
    void dfs(int x,
                    int p) {
      tin[x] = low[x] = ++t;
      int ch = 0;
      for (auto u: G[x]) {
        if (u.first == p) continue;
        if (not ins[u.second]) {
          st.push_back(u.second);
          ins[u.second] = true;
        if (tin[u.first]) {
          low[x] = min(low[x], tin[u.first]);
```

```
continue:
         ++ch; dfs(u.first, x);
         low[x] = min(low[x], low[u.first]);
         if (low[u.first] >= tin[x]) {
           ap[x] = true; ++ecnt;
           while (true) {
             int e = st.back(); st.pop_back();
             bcc[e] = ecnt;
             if (e == u.second) break;
        }
      if (ch == 1 \text{ and } p == x) ap[x] = false;
  public:
    void init(int n_) {
  n = n_, ecnt = 0; st.clear();
      G.clear(); G.resize(n);
      low.clear(); tin.clear();
ap.assign(n, false);
    void add_edge(int u, int v) {
      G[u].emplace_back(v, ecnt);
      G[v].emplace_back(u, ecnt++);
    void solve() {
      ecnt = 0; bcc.resize(t);
      ins.assign(t, false);
for (int i = 0; i < n; ++i)</pre>
         if (low[i] == 0) dfs(i, i);
    int get_id(int x) { return bcc[x];; }
    int count() { return ecnt; }
    bool is_ap(int x) { return ap[x]; }
};
3.4 2-SAT (SCC)
class TwoSat{
  private:
    int n;
    vector<vector<int>> rG,G,sccs;
    vector<int> ord,idx;
    vector<bool> vis,result;
    void dfs(int u){
      vis[u]=true
      for(int v:G[u])
        if(!vis[v]) dfs(v);
      ord.push_back(u);
    void rdfs(int u){
      vis[u]=false;idx[u]=sccs.size()-1;
      sccs.back().push_back(u);
      for(int v:rG[u])
         if(vis[v])rdfs(v);
  public:
    void init(int n_){
      n=n_;G.clear();G.resize(n);
      rG.clear();rG.resize(n)
      sccs.clear();ord.clear();
      idx.resize(n);result.resize(n);
    void add_edge(int u,int v){
      G[u].push_back(v);rG[v].push_back(u);
    void orr(int x,int y){
      if ((x^y)==1)return
      add_edge(x^1,y); add_edge(y^1,x);
    bool solve(){
      vis.clear();vis.resize(n);
      for(int i=0;i<n;++i)</pre>
         if(not vis[i])dfs(i);
      reverse(ord.begin(),ord.end());
      for (int u:ord){
        if(!vis[u])continue;
         sccs.push_back(vector<int>());
         rdfs(u);
```

```
for(int i=0;i<n;i+=2)
    if(idx[i]==idx[i+1])
    return false;

vector<bool> c(sccs.size());
for(size_t i=0;i<sccs.size();++i){
    for(size_t j=0;j<sccs[i].size();++j){
        result[sccs[i][j]]=c[i];
        c[idx[sccs[i][j]^1]]=!c[i];
    }
}
return true;
}
bool get(int x){return result[x];}
inline int get_id(int x){return idx[x];}
inline int count(){return sccs.size();}
} sat2;</pre>
```

3.5 Lowbit Decomposition

```
class LowbitDecomp{
private:
 int time_, chain_, LOG_N;
  vector< vector< int > > G, fa;
  vector< int > tl, tr, chain, chain_st;
  // chain_ : number of chain
  // tl, tr[ u ] : subtree interval in the seq. of u
  // chain_st[ u ] : head of the chain contains u
  // chian[ u ] : chain id of the chain u is on
  inline int lowbit( int x ) {
    return x & ( -x );
 void predfs( int u, int f ) {
  chain[ u ] = 0;
    for ( int v : G[ u ] ) {
      if ( v == f ) continue;
      predfs( v, u );
      if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
        chain[ u ] = chain[ v ];
    if ( not chain[ u ] )
      chain[ u ] = chain_ ++;
  void dfschain( int u, int f ) {
   fa[ u ][ 0 ] = f;

for ( int i = 1 ; i < LOG_N ; ++ i )

fa[ u ][ i ] = fa[ fa[ u ][ i - 1 ] ][ i - 1 ];
    tl[ u ] = time_++
    if ( not chain_st[ chain[ u ] ] )
  chain_st[ chain[ u ] ] = u;
    for ( int v : G[ u ] )
      if ( v != f and chain[ v ] == chain[ u ] )
        dfschain( v, u );
    for ( int v : G[ u ] )
      if ( v != f and chain[ v ] != chain[ u ] )
        dfschain( v, u );
    tr[ u ] = time_;
  inline bool anc( int u, int v ) {
    return tl[ u ] <= tl[ v ] \</pre>
      and tr[ v ] <= tr[ u ];
public:
  inline int lca( int u, int v ) {
    if ( anc( u, v ) ) return u;
    for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
      if ( not anc( fa[ u ][ i ], v ) )
        u = fa[ u ][ i ];
    return fa[ u ][ 0 ];
  void init( int n ) {
    for ( LOG_N = 0 ; ( 1 << LOG_N ) < n ; ++ LOG_N );
    fa.clear();
    fa.resize( n, vector< int >( LOG_N ) );
    G.clear(); G.resize( n );
    tl.clear(); tl.resize( n );
    tr.clear(); tr.resize( n );
    chain.clear(); chain.resize( n );
    chain_st.clear(); chain_st.resize( n );
  void add_edge( int u , int v ) {
```

```
// 1-base
    G[ u ].push_back( v );
    G[ v ].push_back( u );
  void decompose(){
    chain_ = 1;
    predfs( 1, 1 );
    time_{-} = 0;
    dfschain( 1, 1 );
  PII get_inter( int u ) { return {tl[ u ], tr[ u ]}; }
  vector< PII > get_path( int u , int v ){
    vector< PII > res;
    int g = lca( u, v );
while ( chain[ u ] != chain[ g ] ) {
      int s = chain_st[ chain[ u ] ];
      res.emplace_back( tl[ s ], tl[ u ] + 1 );
      u = fa[ s ][ 0 ];
    res.emplace_back( tl[ g ], tl[ u ] + 1 );
while ( chain[ v ] != chain[ g ] ) {
      int s = chain_st[ chain[ v ] ];
      res.emplace_back( tl[ s ], tl[ v ] + 1 );
      v = fa[ s ][ 0 ];
    res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
    return res;
    /* res : list of intervals from u to v
     \star ( note only nodes work, not edge )
     * usage :
     * vector< PII >& path = tree.get_path( u , v )
     * for( auto [ 1, r ] : path ) {
         0-base [ 1, r )
} tree;
```

3.6 MaxClique

```
// contain a self loop u to u, than u won't in clique
template < size_t MAXN >
class MaxClique{
private:
  using bits = bitset< MAXN >;
  bits popped, G[ MAXN ], ans;
  size_t deg[ MAXN ], deo[ MAXN ], n;
  void sort_by_degree() {
    popped.reset();
    for ( size_t i = 0 ; i < n ; ++ i )
    deg[ i ] = G[ i ].count();
for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
         size_t mi = MAXN, id = 0;
         for ( size_t j = 0 ; j < n ; ++ j )
    if ( not popped[ j ] and deg[ j ] < mi )</pre>
                  mi = deg[ id = j ];
         popped[ deo[ i ] = id ] = 1;
for( size_t u = G[ i ]._Find_first()
           u < n ; u = G[ i ]._Find_next( u ) )
             -- deg[ u ];
  void BK( bits R, bits P, bits X ) {
    if (R.count()+P.count() <= ans.count()) return;</pre>
    if ( not P.count() and not X.count() ) {
      if ( R.count() > ans.count() ) ans = R;
      return;
    /* greedily chosse max degree as pivot
    bits cur = P | X; size_t pivot = 0, sz = 0;
    for ( size_t u = cur._Find_first() ;
      u < n ; u = cur._Find_next( u )</pre>
         if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
    cur = P & ( ~G[ pivot ] );
    */ // or simply choose first
    bits cur = P & (~G[ ( P | X )._Find_first() ]);
    for ( size_t u = cur._Find_first() ;
      u < n ; u = cur._Find_next( u ) ) {</pre>
      if ( R[ u ] ) continue;
      R[u] = 1;
      BK( R, P & G[ u ], X & G[ u ] );
```

```
R[u] = P[u] = 0, X[u] = 1;
     }
public:
  void init( size_t n_ ) {
     n = n_{-};
     for ( size_t i = 0 ; i < n ; ++ i )
  G[ i ].reset();</pre>
     ans.reset();
   void add_edges( int u, bits S ) { G[ u ] = S; }
  void add_edge( int u, int v ) {
     G[u][v] = G[v][u] = 1;
   int solve() {
     sort_by_degree(); // or simply iota( deo... )
     for ( size_t i = 0 ; i < n ; ++ i )
  deg[ i ] = G[ i ].count();</pre>
     bits pob, nob = 0; pob.set();
     for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
     for ( size_t i = 0 ; i < n ; ++ i ) {
       size_t v = deo[ i ];
       bits tmp; tmp[ v ] = 1;

BK( tmp, pob & G[ v ], nob & G[ v ] );

pob[ v ] = 0, nob[ v ] = 1;
     return static_cast< int >( ans.count() );
|};
```

3.7 Virtural Tree

```
inline bool cmp(const int &i, const int &j) {
  return dfn[i] < dfn[j];</pre>
void build(int vectrices[], int k) {
 static int stk[MAX_N];
  sort(vectrices, vectrices + k, cmp);
  stk[sz++] = 0;
 for (int i = 0; i < k; ++i) {
  int u = vectrices[i], lca = LCA(u, stk[sz - 1]);</pre>
    if (lca == stk[sz - 1]) stk[sz++] = u;
    else {
      while (sz >= 2 && dep[stk[sz - 2]] >= dep[lca]) {
        addEdge(stk[sz - 2], stk[sz - 1]);
        sz--;
      if (stk[sz - 1] != lca) {
        addEdge(lca, stk[--sz]);
        stk[sz++] = lca, vectrices[cnt++] = lca;
      stk[sz++] = u;
  for (int i = 0; i < sz - 1; ++i)
    addEdge(stk[i], stk[i + 1]);
```

3.8 Tree Hashing

```
uint64_t hsah( int u, int f ) {
    uint64_t r = 127;
    for ( int v : G[ u ] ) {
        if ( v == f ) continue;
        uint64_t hh = hsah( v, u );
        r = r + ( hh * hh ) % mod;
    }
    return r;
}
```

3.9 Minimum Mean Cycle

```
/* minimum mean cycle O(VE) */
struct MMC{
    #define FZ(n) memset((n),0,sizeof(n))
    #define E 101010
    #define V 1021
#define inf 1e9
    struct Edge { int v,u; double c; };
    int n, m, prv[V][V], prve[V][V], vst[V];
```

```
Edge e[E];
  vector<int> edgeID, cycle, rho;
  double d[V][V];
  void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
  void add_edge( int vi , int ui , double ci )
  { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
    for(int i=0; i<n; i++) d[0][i]=0;
    for(int i=0; i<n; i++) {</pre>
      fill(d[i+1], d[i+1]+n, inf);
for(int j=0; j<m; j++) {
         int v = e[j].v, u = e[j].u;
         if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
           d[i+1][u] = d[i][v]+e[j].c;
           prv[i+1][u] = v;
           prve[i+1][u] = j;
    }
  double solve(){
    // returns inf if no cycle, mmc otherwise
    double mmc=inf;
    int st = -1;
    bellman_ford();
    for(int i=0; i<n; i++) {</pre>
      double avg=-inf;
       for(int k=0; k<n; k++) {</pre>
         if(d[n][i]<inf-eps)</pre>
           avg=max(avg,(d[n][i]-d[k][i])/(n-k));\\
         else avg=max(avg,inf);
      if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
    FZ(vst);edgeID.clear();cycle.clear();rho.clear();
    for (int i=n; !vst[st]; st=prv[i--][st]) {
      vst[st]++;
      edgeID.PB(prve[i][st]);
      rho.PB(st);
    while (vst[st] != 2) {
      int v = rho.back(); rho.pop_back();
      cycle.PB(v);
      vst[v]++;
    reverse(ALL(edgeID));
    edgeID.resize(SZ(cycle));
    return mmc;
} mmc;
```

3.10 Mo's Algorithm on Tree

```
int n, q, nxt[ N ], to[ N ], hd[ N ];
struct Que{
  int u, v, id;
} que[ N ];
void init() {
  cin >> n >> q;
  for ( int i = 1 ; i < n ; ++ i ) {
  int u, v; cin >> u >> v;
     nxt[ i << 1 | 0 ] = hd[ u ];</pre>
     to[ i << 1 | 0 ] = v;
hd[ u ] = i << 1 | 0;
     nxt[ i << 1 | 1 ] = hd[ v ];
     to[ i << 1 | 1 ] = u;
     hd[ v ] = i << 1 | 1;
  for ( int i = 0 ; i < q ; ++ i ) {
  cin >> que[ i ].u >> que[ i ].v; que[ i ].id = i;
int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
void dfs( int u, int f ) {
  dfn[ u ] = dfn_++; int saved_rbp = stk_;
for ( int v_ = hd[ u ] ; v_ ; v_ = nxt[ v_ ] ) {
     if ( to[ v_ ] == f ) continue;
     dfs( to[ v_ ], u );
if ( stk_ - saved_rbp < SQRT_N ) continue;</pre>
     for ( ++ block_ ; stk_ != saved_rbp ; )
```

```
block_id[ stk[ -- stk_ ] ] = block_;
  stk[ stk_ ++ ] = u;
bool inPath[ N ];
void Diff( int u )
  if ( inPath[ u ] ^= 1 )
    // remove this edge
  else
    // add this edge
void traverse( int& origin_u, int u ) {
  for ( int g = lca( origin_u, u ) ;
    origin_u != g ; origin_u = parent_of[ origin_u ] )
      Diff( origin_u );
  for (int v = u; v != origin_u; v = parent_of[v])
    Diff( v );
  origin_u = u;
void solve() {
 dfs( 1, 1 );
  while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
  sort( que, que + q, [](const Que& x, const Que& y) {
  return tie( block_id[ x.u ], dfn[ x.v ] )
             < tie( block_id[ y.u ], dfn[ y.v ] );
  } );
  int U = 1, V = 1;
  for ( int i = 0 ; i < q ; ++ i ) {</pre>
    pass( U, que[ i ].u );
pass( V, que[ i ].v );
    // we could get our answer of que[ i ].id
}
/*
Method 2:
dfs u:
  push u
  iterate subtree
 push u
Let P = LCA(u, v), and St(u) <= St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
*/
```

3.11 Minimum Steiner Tree

```
// Minimum Steiner Tree
// 0(V 3^T + V^2 2^T)
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
  int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
  void init( int _n ){
    n = _n;
for( int i = 0 ; i < n ; i ++ ){</pre>
       for( int j = 0 ; j < n ; j ++ ) {
   dst[ i ][ j ] = INF;
   dst[ i ][ i ] = 0;
     }
  void add_edge( int ui , int vi , int wi ){
     dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
     dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
  void shortest_path(){
     for( int k = 0 ; k < n ; k ++ )
  for( int i = 0 ; i < n ; i ++ )</pre>
          for( int j = 0 ; j < n ; j ++ )
  dst[ i ][ j ] = min( dst[ i ][ j ],</pre>
                    dst[i][k]+dst[k][j]);
  int solve( const vector<int>& ter ){
     int t = (int)ter.size();
     for( int i = 0 ; i < ( 1 << t ) ; i ++ )
     for( int j = 0 ; j < n ; j ++ )
   dp[ i ][ j ] = INF;
for( int i = 0 ; i < n ; i ++ )</pre>
        dp[0][i] = 0;
     for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
       if( msk == ( msk & (-msk) ) ){
```

```
int who = __lg( msk );
for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];</pre>
           continue:
        for( int i = 0 ; i < n ; i ++ )
           for( int submsk = ( msk - 1 ) & msk ; submsk ;
    submsk = ( submsk - 1 ) & msk )
                 dp[ msk ][ i ] = min( dp[ msk ][ i ],
                                        dp[ submsk ][ i ] +
dp[ msk ^ submsk ][ i ] );
        for( int i = 0 ; i < n ; i ++ ){</pre>
            tdst[ i ] = INF;
           for( int j = 0 ; j < n ; j ++ )
  tdst[ i ] = min( tdst[ i ],</pre>
                              dp[ msk ][ j ] + dst[ j ][ i ] );
        for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = tdst[ i ];</pre>
      int ans = INF;
      for( int i = 0 ; i < n ; i ++ )</pre>
        ans = min(ans, dp[(1 << t) - 1][i]);
      return ans:
} solver;
```

3.12 Directed Minimum Spanning Tree

```
\begin{tabular}{lll} \textbf{template} & \textbf{<typename} & \textbf{T>} & \textbf{struct} & \textbf{DMST} & \textbf{&} \\ \end{tabular}
  T g[maxn][maxn], fw[maxn];
  int n, fr[maxn];
  bool vis[maxn], inc[maxn];
  void clear() {
     for(int i = 0; i < maxn; ++i) {</pre>
       for(int j = 0; j < maxn; ++j) g[i][j] = inf;
       vis[i] = inc[i] = false;
    }
  void addEdge(int u,int v,T w){g[u][v]=min(g[u][v],w)
  T operator()(int root, int _n) {
     n = _n; T ans = 0;
     if (dfs(root) != n) return -1;
     while (true) {
  for(int i = 1;i <= n;++i) fw[i] = inf, fr[i] = i;</pre>
       for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
          for (int j = 1; j <= n; ++j) {
  if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
              fw[i] = g[j][i]; fr[i] = j;
            }
         }
       int x = -1;
       for(int i = 1;i <= n;++i)if(i != root && !inc[i])</pre>
         int j = i, c = 0;
         while(j!=root && fr[j]!=i && c<=n) ++c, j=fr[j</pre>
          if (j == root || c > n) continue;
         else { x = i; break; }
       if (!~x) {
          for (int i = 1; i <= n; ++i)
            if (i != root && !inc[i]) ans += fw[i];
          return ans:
       int y = x;
       for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
       do {
         ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true
       } while (y != x);
       inc[x] = false;
       for (int k = 1; k <= n; ++k) if (vis[k])</pre>
          for (int j = 1; j <= n; ++j) if (!vis[j]) {
            if^{(g[x][j])} = g[k][j] = g[x][j] = g[k][j];
if^{(g[j][k])} < inf^{(g[j][k])} = g[k][j][x]
               g[j][x] = g[j][k] - fw[k];
         }
```

```
if (vl[x] = true, fl[x] != -1)
                                                                    return vr[qu[qr++] = fl[x]] = true;
    return ans;
                                                                  while (x != -1) swap(x, fr[fl[x] = pre[x]]);
                                                                  return false:
  int dfs(int now) {
    int r = 1; vis[now] = true;
    for (int i = 1; i <= n; ++i)
                                                               void bfs(int s) {
      fill(slk.begin(), slk.end(), INF);
fill(vl.begin(), vl.end(), false);
                                                                  fill(vr.begin(), vr.end(), false);
};
                                                                  ql = qr = 0;
                                                                  qu[qr++] = s;
                                                                  vr[s] = true;
3.13 Dominator Tree
                                                                  while (true) {
                                                                    11d d;
namespace dominator {
                                                                    while (ql < qr) {</pre>
vector<int> g[maxn], r[maxn], rdom[maxn];
                                                                      for (int x = 0, y = qu[ql++]; x < n; ++x)
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
                                                                        if(!v1[x]\&\&s1k[x]>=(d=h1[x]+hr[y]-w[x][y])){
                                                                          if (pre[x] = y, d) slk[x] = d;
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
                                                                          else if (!check(x)) return;
  // vertices are numbered from 0 to n - 1 \,
                                                                        }
  fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
                                                                     }
  fill(fa, fa + n, -1); fill(val, val + n, -1);
  fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
                                                                    d = INF;
  fill(dom, dom + n, -1); tk = 0;
                                                                    for (int x = 0; x < n; ++x)
 for (int i = 0; i < n; ++i) {
  g[i].clear(); r[i].clear(); rdom[i].clear();</pre>
                                                                      if (!vl[x] && d > slk[x]) d = slk[x];
                                                                    for (int x = 0; x < n; ++x) {
                                                                      if (vl[x]) hl[x] += d;
                                                                      else slk[x] -= d;
void add_edge(int x, int y) { g[x].push_back(y); }
                                                                      if (vr[x]) hr[x] -= d;
void dfs(int x) {
  rev[dfn[x] = tk] = x;
                                                                    for (int x = 0; x < n; ++x)
  fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
                                                                      if (!v1[x] && !slk[x] && !check(x)) return;
  for (int u : g[x]) {
    if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
                                                               }
    r[dfn[u]].push_back(dfn[x]);
                                                             public:
  }
                                                               void init( int n_ ) {
                                                                  n = n_; qu.resize(n);
void merge(int x, int y) { fa[x] = y; }
                                                                  fl.clear(); fl.resize(n, -1);
int find(int x, int c = 0)
                                                                 fr.clear(); fr.resize(n, -1);
hr.clear(); hr.resize(n); hl.resize(n);
  if (fa[x] == x) return c ? -1 : x;
  int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
                                                                  w.clear(); w.resize(n, vector<lld>(n));
                                                                  slk.resize(n); pre.resize(n);
  if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
                                                                  vl.resize(n); vr.resize(n);
  fa[x] = p;
return c ? p : val[x];
                                                               void set_edge( int u, int v, lld x ) {w[u][v] = x;}
                                                               1ld solve() {
vector<int> build(int s, int n) {
                                                                  for (int i = 0; i < n; ++i)
// return the father of each node in the dominator tree
                                                                   hl[i] = *max_element(w[i].begin(), w[i].end());
// p[i] = -2 if i is unreachable from s
                                                                  for (int i = 0; i < n; ++i) bfs(i);</pre>
  dfs(s);
                                                                  11d res = 0;
  for (int i = tk - 1; i >= 0; --i) {
                                                                  for (int i = 0; i < n; ++i) res += w[i][f1[i]];</pre>
    for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)])
                                                                  return res;
```

} km;

}

vector<int> p(n, -2); p[s] = -1; for (int i = 1; i < tk; ++i)</pre> vector<int> X[N], Y[N]; int fX[N], fY[N], n; bitset<N> walked:

4 Matching & Flow

if (i) rdom[sdom[i]].push_back(i);

if (sdom[p] == i) dom[u] = i;

for (int &u : rdom[i]) {
 int p = find(u);

else dom[u] = p;

if (i) merge(i, rp[i]);

4.1 Kuhn Munkres

return p;

```
class KM {
private:
    static constexpr lld INF = 1LL << 60;
    vector<lld> hl,hr,slk;
    vector<int> fl,fr,pre,qu;
    vector<vector<lld> w;
    vector<bool> vl,vr;
    int n, ql, qr;
    bool check(int x) {
```

if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];

for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>

```
Bipartite Matching
class BipartiteMatching{
private:
  vector<int> X[N], Y[N];
  int fX[N], fY[N], n;
bitset<N> walked;
  bool dfs(int x){
    for(auto i:X[x]){
      if(walked[i])continue;
      walked[i]=1;
      if(fY[i]==-1||dfs(fY[i])){
        fY[i]=x;fX[x]=i;
        return 1;
      }
    return 0;
public:
  void init(int _n){
    n=_n; walked.reset();
    for(int i=0;i<n;i++){</pre>
      X[i].clear();Y[i].clear();
      fX[i]=fY[i]=-1;
```

```
void add_edge(int x, int y){
  X[x].push_back(y); Y[y].push_back(y);
  int solve(){
     int cnt = 0;
     for(int i=0;i<n;i++){</pre>
       walked.reset();
       if(dfs(i)) cnt++;
     // return how many pair matched
     return cnt:
  }
};
```

General Graph Matching

```
const int N = 514, E = (2e5) * 2;
struct Graph{
  int to[E],bro[E],head[N],e;
  int lnk[N], vis[N], stp, n;
  void init( int _n ){
    stp = 0; e = 1; n = _n;
for( int i = 0 ; i <= n ; i ++ )
      head[i] = lnk[i] = vis[i] = 0;
  void add_edge(int u,int v){
    // 1-base
    to[e]=v,bro[e]=head[u],head[u]=e++;
    to[e]=u, bro[e]=head[v], head[v]=e++;
  bool dfs(int x){
    vis[x]=stp;
    for(int i=head[x];i;i=bro[i]){
      int v=to[i];
      if(!lnk[v]){
        lnk[x]=v, lnk[v]=x;
        return true;
      }else if(vis[lnk[v]]<stp){</pre>
         int w=lnk[v];
        lnk[x]=v, lnk[v]=x, lnk[w]=0;
         if(dfs(w)) return true
        lnk[w]=v, lnk[v]=w, lnk[x]=0;
      }
    }
    return false;
  int solve(){
    int ans = 0;
    for(int i=1;i<=n;i++)</pre>
      if(not lnk[i]){
        stp++; ans += dfs(i);
    return ans;
} graph;
```

Minimum Weight Matching (Clique 4.4 version)

```
struct Graph {
  // 0-base (Perfect Match)
  int n, edge[MXN][MXN];
  int match[MXN], dis[MXN], onstk[MXN];
  vector<int> stk;
  void init(int _n) {
   n = _n;
for (int i=0; i<n; i++)</pre>
      for (int j=0; j<n; j++)</pre>
        edge[i][j] = 0;
  void set_edge(int u, int v, int w) {
    edge[u][v] = edge[v][u] = w;
  bool SPFA(int u){
    if (onstk[u]) return true;
    stk.PB(u);
    onstk[u] = 1;
    for (int v=0; v<n; v++){
      if (u != v && match[u] != v && !onstk[v]){
        int m = match[v];
```

```
if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
           dis[m] = dis[u] - edge[v][m] + edge[u][v];
           onstk[v] = 1;
           stk.PB(v);
           if (SPFA(m)) return true;
           stk.pop_back();
           onstk[v] = 0;
       }
     onstk[u] = 0;
     stk.pop_back();
     return false;
   int solve() {
     // find a match
     for (int i=0; i<n; i+=2){
       match[i] = i+1;
       match[i+1] = i;
     while (true){
       int found = 0;
       for (int i=0; i<n; i++)</pre>
         dis[i] = onstk[i] = 0;
       for (int i=0; i<n; i++){</pre>
         stk.clear()
         if (!onstk[i] && SPFA(i)){
           found = 1
           while (SZ(stk)>=2){
             int u = stk.back(); stk.pop_back();
             int v = stk.back(); stk.pop_back();
             match[u] = v;
             match[v] = u;
         }
       if (!found) break;
     int ret = 0;
     for (int i=0; i<n; i++)</pre>
       ret += edge[i][match[i]];
     return ret>>1;
  }
} graph;
```

4.5 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\it S}$ and sink ${\it T}$.
 - 2. For each edge (x,y,l,u), connect x o y with capacity
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow
 - from s to t is the answer. - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f^\prime is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corre-
- sponds to the flow of edge $\it e$ on the graph. ullet Construct minimum vertex cover from maximum matching M on
 - $\ \, \text{bipartite graph}\ (X,Y)$ 1. Redirect every edge: $y \rightarrow x$ if $(x,y) \in M$, $x \rightarrow y$ otherwise.

 - 2. DFS from unmatched vertices in X. 3. $x \in X$ is chosen iff x is unvisited. 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\it S}$ and sink ${\it T}$
 - 2. For each edge (x,y,c), connect $x \to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y \to x$ with (cost,cap)=(c,1)
 - 3. For each edge with c < 0, sum these cost as K, then
 - increase d(y) by 1, decrease d(x) by 1 4. For each vertex v with d(v)>0, connect $S\to v$ with $(\cos t, cap) = (0, d(v))$

```
5. For each vertex v with d(v) < 0, connect v 	o T with
    (cost, cap) = (0, -d(v)) 6. Flow from S to T, the answer is the cost of the flow
· Maximum density induced subgraph
    1. Binary search on answer, suppose we're checking answer
    2. Construct a max flow model. let K be the sum of all
```

- weights 3. Connect source s o v, $v \in G$ with capacity K
- 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity \boldsymbol{w}
- 5. For $v\in G$, connect it with sink $v\to t$ with capacity $K+2T-(\sum_{e\in E(v)}w(e))-2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect u' o v'with weight w(u,v). 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost
 - of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise,
 - create edge (v,t) with capacity $-p_v$. 2. Create edge (u,v) with capacity w with w being the cost
 - of choosing u without choosing $v\,.$ 3. The mincut is equivalent to the maximum profit of a subset of projects.
- · 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y)with capacity $c_y\,.$
- 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4.6 Dinic

class Dinic{

```
private:
 using CapT = int64_t;
  struct Edge{
    int to, rev;
   CapT cap;
 int n, st, ed;
 vector<vector<Edge>> G;
  vector<int> lv, idx;
 bool BFS(){
    fill(lv.begin(), lv.end(), -1);
    queue<int> bfs;
    bfs.push(st);
    lv[st] = 0;
    while(!bfs.empty()){
      int u = bfs.front(); bfs.pop();
      for(auto e: G[u]){
        if(e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
        lv[e.to] = lv[u] + 1;
        bfs.push(e.to);
      }
    return (lv[ed]!=-1);
 CapT DFS(int u, CapT f){
    if(u == ed) return f;
    CapT ret = 0;
    for(int& i = idx[u]; i < (int)G[u].size(); ++i){</pre>
      auto& e = G[u][i]
      if(e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;
      CapT nf = DFS(e.to, min(f, e.cap));
      ret += nf; e.cap -= nf; f -= nf;
      G[e.to][e.rev].cap += nf;
      if(f == 0) return ret;
    if(ret == 0) lv[u] = -1;
    return ret;
public:
 void init(int n_, int st_, int ed_){
   n = n_, st = st_, ed = ed_;
    G.resize(n); lv.resize(n);
```

```
fill(G.begin(), G.end(), vector<Edge>());
  void add_edge(int u, int v, CapT c){
    G[u].push_back({v, (int)G[v].size(), c});
G[v].push_back({u, ((int)G[u].size())-1, 0});
  CapT max_flow(){
    CapT ret = 0;
    while(BFS()){
       idx.assign(n, 0);
       CapT f = DFS(st, numeric_limits<CapT>::max());
       ret += f;
       if(f == 0) break;
    return ret;
} flow;
```

4.7 Minimum Cost Maximum Flow

```
class MiniCostMaxiFlow{
  using CapT = int;
  using WeiT = int64_t;
  using PCW = pair<CapT,WeiT>;
  static constexpr CapT INF_CAP = 1 << 30;</pre>
  static constexpr WeiT INF_WEI = 1LL<<60;</pre>
private:
  struct Edge{
    int to, back;
    WeiT wei;
    CapT cap;
    Edge() {}
    Edge(int a,int b,WeiT c,CapT d):
      to(a),back(b),wei(c),cap(d)
    {}
  int ori, edd;
  vector<vector<Edge>> G;
  vector<int> fa, wh;
  vector<bool> ing;
  vector<WeiT> dis;
  PCW SPFA(){
    fill(inq.begin(),inq.end(),false);
    fill(dis.begin(), dis.end(), INF_WEI);
    queue<int> qq; qq.push(ori);
    dis[ori]=0;
    while(!qq.empty()){
      int u=qq.front();qq.pop();
      inq[u] = 0;
      for(int i=0;i<SZ(G[u]);++i){</pre>
        Edge e=G[u][i];
        int v=e.to;
        WeiT d=e.wei;
        if(e.cap <= 0 \mid |dis[v] <= dis[u] + d)
           continue
        dis[v]=dis[u]+d:
        fa[v]=u,wh[v]=i;
        if(inq[v]) continue;
        qq.push(v);
        inq[v]=1;
      }
    if(dis[edd]==INF_WEI)
      return {-1,-1};
    CapT mw=INF_CAP;
    for(int i=edd;i!=ori;i=fa[i])
      mw=min(mw,G[fa[i]][wh[i]].cap);
    for (int i=edd;i!=ori;i=fa[i]){
      auto &eg=G[fa[i]][wh[i]];
      eq.cap-=mw:
      G[eg.to][eg.back].cap+=mw;
    return {mw,dis[edd]};
  }
public:
  void init(int a,int b,int n){
    ori=a,edd=b;
    G.clear();G.resize(n);
    fa.resize(n);wh.resize(n);
    inq.resize(n); dis.resize(n);
```

```
void add_edge(int st,int ed,WeiT w,CapT c){
    G[st].emplace_back(ed,SZ(G[ed]),w,c);
    G[ed].emplace_back(st,SZ(G[st])-1,-w,0);
 PCW solve(){
    /* might modify to
    cc += ret.first * ret.second
   ww += ret.first * ret.second
    */
   CapT cc=0; WeiT ww=0;
   while(true){
      PCW ret=SPFA();
      if(ret.first==-1) break;
      cc+=ret.first;
      ww+=ret.second;
    return {cc,ww};
} mcmf;
```

4.8 Global Min-Cut

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
    w[x][y] += c;
    w[y][x] += c;
pair<int, int> phase(int n) {
    memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
    int s = -1, t = -1;
    while (true) {
         int c = -1;
         for (int i = 0; i < n; ++i) {
              if (del[i] || v[i]) continue;
              if (c == -1 \mid \mid g[i] > g[c]) c = i;
         if (c == -1) break;
         v[c] = true;
         s = t, t = c;
         for (int i = 0; i < n; ++i) {
              if (del[i] || v[i]) continue;
              g[i] += w[c][i];
     return make_pair(s, t);
}
int mincut(int n) {
    int cut = 1e9;
    memset(del, false, sizeof(del));
for (int i = 0; i < n - 1; ++i) {</pre>
         int s, t; tie(s, t) = phase(n);
         del[t] = true;
         cut = min(cut, g[t]);
for (int j = 0; j < n; ++j) {
    w[s][j] += w[t][j];</pre>
              w[j][s] += w[j][t];
         }
     return cut;
}
```

5 Math

5.1 Prime Table

```
\begin{array}{c} 1002939109, 1020288887, 1028798297, 1038684299, \\ 1041211027, 1051762951, 1058585963, 1063020809, \\ 1147930723, 1172520109, 1183835981, 1187659051, \\ 1241251303, 1247184097, 1255940849, 1272759031, \\ 1287027493, 1288511629, 1294632499, 1312650799, \\ 1868732623, 1884198443, 1884616807, 1885059541, \\ 1909942399, 1914471137, 1923951707, 1925453197, \\ 1979612177, 1980446837, 1989761941, 2007826547, \\ 2008033571, 2011186739, 2039465081, 2039728567, \\ 2093735719, 2116097521, 2123852629, 2140170259, \end{array}
```

```
\begin{array}{c} 3148478261, 3153064147, 3176351071, 3187523093, \\ 3196772239, 3201312913, 3203063977, 3204840059, \\ 3210224309, 3213032591, 3217689851, 3218469083, \\ 3219857533, 3231880427, 3235951699, 3273767923, \\ 3276188869, 3277183181, 3282463507, 3285553889, \\ 3319309027, 3327005333, 3327574903, 3341387953, \\ 3373293941, 3380077549, 3380892997, 3381118801 \end{array}
```

5.2 $\lfloor \frac{n}{i} \rfloor$ Enumeration

```
T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_{i+1} \rfloor} \rfloor}
```

5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b==(x/y)*a;
}
```

5.4 Pollard Rho

```
// does not work when n is prime
// return any non-trivial factor
llu pollard_rho(llu n){
  static auto f=[](llu x,llu k,llu m){
    return add(k,mul(x,x,m),m);
  if (!(n&1)) return 2;
  mt19937 rnd(120821011);
  while(true){
    llu y=2, yy=y, x=rnd()%n, t=1;
    for(llu sz=2;t==1;sz<<=1) {</pre>
      for(llu i=0;i<sz;++i){</pre>
        if(t!=1)break;
        yy=f(yy,x,n);
        t=gcd(yy>y?yy-y:y-yy,n);
      y=yy;
    if(t!=1&&t!=n) return t;
}
```

5.5 Pi Count (Linear Sieve)

```
static constexpr int N = 1000000 + 5;
11d pi[N];
vector<int>
             primes;
bool sieved[N];
11d cube_root(11d x){
  1ld_s=cbrt(x-static_cast<long double>(0.1));
  while(s*s*s <= x) ++s;
  return s-1;
1ld square_root(lld x){
  1ld s=sqrt(x-static_cast<long double>(0.1));
  while(s*s <= x) ++s;
  return s-1;
void init(){
  primes.reserve(N)
  primes.push_back(1);
  for(int i=2;i<N;i++) {</pre>
    if(!sieved[i]) primes.push_back(i);
    pi[i] = !sievéd[i] + pi[i-1];
    for(int p: primes) if(p > 1) {
  if(p * i >= N) break;
  sieved[p * i] = true;
}
       if(p % i == 0) break;
  }
ild phi(lld m, lld n) {
   static constexpr int MM = 80000, NN = 500;
  static lld val[MM][NN];
  if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
  if(n == 0) return m;
  if(primes[n] >= m) return 1;
  1ld ret = phi(m,n-1)-phi(m/primes[n],n-1);
  if(m < MM\&n < NN) val[m][n] = ret+1;
```

return ret;

```
}
lld pi_count(lld);
lld P2(lld m, lld n) {
    lld sm = square_root(m), ret = 0;
    for(lld i = n+1;primes[i]<=sm;i++)
        ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
    return ret;
}
lld pi_count(lld m) {
    if(m < N) return pi[m];
    lld n = pi_count(cube_root(m));
    return phi(m, n) + n - 1 - P2(m, n);
}
</pre>
```

5.6 Range Sieve

5.7 Miller Rabin

```
bool isprime(llu x){
 static llu magic[]={2,325,9375,28178,\
                    450775,9780504,1795265022};
 static auto witn=[](llu a,llu u,llu n,int t){
    a = mpow(a,u,n)
    if (!a)return 0;
   while(t--){
      1lu a2=mul(a,a,n);
      if(a2==1 && a!=1 && a!=n-1)
        return 1;
      a = a2:
    }
    return a!=1;
 if(x<2)return 0;</pre>
  if(!(x&1))return x==2;
 llu x1=x-1;int t=0;
 while(!(x1&1))x1>>=1,t++;
  for(llu m:magic)if(witn(m,x1,x,t))return 0;
 return 1:
```

5.8 Inverse Element

```
// x's inverse mod k
long long GetInv(long long x, long long k){
    // k is prime: euler_(k)=k-1
    return qPow(x, euler_phi(k)-1);
}
// if you need [1, x] (most use: [1, k-1]
void solve(int x, long long k){
    inv[1] = 1;
    for(int i=2;i<x;i++)
        inv[i] = ((long long)(k - k/i) * inv[k % i]) % k;
}</pre>
```

5.9 Euler Phi Function

```
extended euler:
a^b mod p
if gcd(a, p)==1: a^(b%phi(p))
```

```
elif b < phi(p): a^b mod p
   else a^(b%phi(p) + phi(p))
lld euler_phi(int x){
  11d r=1;
  for(int i=2;i*i<=x;++i){</pre>
    if(x%i==0){
      x/=i; r*=(i-1);
      while(x%i==0){
        x/=i; r*=i;
    }
  if(x>1) r*=x-1;
  return r;
vector<int> primes;
bool notprime[N];
11d phi[N];
void euler_sieve(int n){
  for(int i=2;i<n;i++){</pre>
    if(!notprime[i]){
      primes.push_back(i); phi[i] = i-1;
    for(auto j: primes){
  if(i*j >= n) break;
      notprime[i*j] = true;
      phi[i*j] = phi[i] * phi[j];
       if(i % j == 0){
        phi[i*j] = phi[i] * j;
        break:
      }
  }
}
```

5.10 Gauss Elimination

```
void gauss(vector<vector<double>> &d) {
   int n = d.size(), m = d[0].size();
   for (int i = 0; i < m; ++i) {
      int p = -1;
      for (int j = i; j < n; ++j) {
        if (fabs(d[j][i]) < eps) continue;
        if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
      }
   if (p == -1) continue;
   for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
   for (int j = 0; j < n; ++j) {
      if (i == j) continue;
      double z = d[j][i] / d[i][i];
      for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
   }
}</pre>
```

5.11 Fast Fourier Transform

```
polynomial multiply:
   DFT(a, len); DFT(b, len);
   for(int i=0;i<len;i++) c[i] = a[i]*b[i];
   iDFT(c, len);
   (len must be 2^k and = 2^k(max(a, b)))
   Hand written Cplx would be 2x faster
Cplx omega[2][N];
void init_omega(int n)
  static constexpr llf PI=acos(-1);
  const llf arg=(PI+PI)/n;
  for(int i=0;i<n;++i)</pre>
    omega[0][i]={cos(arg*i),sin(arg*i)};
  for(int i=0;i<n;++i)</pre>
    omega[1][i]=conj(omega[0][i]);
void tran(Cplx arr[],int n,Cplx omg[]) {
  for(int i=0, j=0;i<n;++i){</pre>
    if(i>j)swap(arr[i],arr[j]);
    for(int l=n>>1;(j^=1)<1;l>>=1);
  for (int 1=2;1<=n;1<<=1){
```

```
int m=l>>1;
  for(auto p=arr;p!=arr+n;p+=l){
    for(int i=0;i<m;++i){
        Cplx t=omg[n/l*i]*p[m+i];
        p[m+i]=p[i]-t; p[i]+=t;
    }
  }
}

void DFT(Cplx arr[],int n){tran(arr,n,omega[0]);}

void iDFT(Cplx arr[],int n){
  tran(arr,n,omega[1]);
  for(int i=0;i<n;++i) arr[i]/=n;
}</pre>
```

5.12 High Speed Linear Recurrence

```
#define mod 998244353
const int N=1000010;
int n,k,m,f[N],h[N],a[N],b[N],ib[N];
int pw(int x,int y){
  int re=1;
  if(y<0)y+=mod-1;
  while(y)
    if(y&1)re=(11)re*x%mod;
    y>=1; x=(11)x*x%mod;
  return re;
void inc(int&x,int y){x+=y;if(x>=mod)x-=mod;}
namespace poly{
  const int G=3;
  int rev[N],L;
  void ntt(int*A,int len,int f){
    for(L=0;(1<<L)<len;++L);</pre>
    for(int i=0;i<len;++i){</pre>
      rev[i]=(rev[i>>1]>>1)|((i&1)<<(L-1));
      if(i<rev[i])swap(A[i],A[rev[i]]);</pre>
    for(int i=1;i<len;i<<=1){</pre>
      int wn=pw(G, f*(mod-1)/(i<<1));</pre>
      for(int j=0;j<len;j+=i<<1){</pre>
        int w=1:
         for(int k=0;k<i;++k,w=(11)w*wn%mod){</pre>
           int x=A[j+k],y=(11)w*A[j+k+i]%mod
           A[j+k]=(x+y)\text{%mod}, A[j+k+i]=(x-y+mod)\text{%mod};
      }
    if(!~f){
      int iv=pw(len,mod-2);
      for(int i=0;i<len;++i)A[i]=(11)A[i]*iv%mod;</pre>
  void cls(int*A,int l,int r){
    for(int i=1;i<r;++i)A[i]=0;}</pre>
  void cpy(int*A,int*B,int 1){
    for(int i=0;i<1;++i)A[i]=B[i];}</pre>
  void inv(int*A,int*B,int 1){
    if(l==1){B[0]=pw(A[0],mod-2);return;}
    static int t[N];
    int len=l<<1:
    inv(A,B,1>>1);
    cpy(t, A, 1); cls(t, 1, len);
    ntt(t,len,1);ntt(B,len,1);
    for(int i=0;i<len;++i)</pre>
      B[i]=(11)B[i]*(2-(11)t[i]*B[i]%mod+mod)%mod;
    ntt(B, len, -1); cls(B, 1, len);
  void pmod(int*A)
    static int t[N];
    int l=k+1,len=1;while(len<=(k<<1))len<<=1;</pre>
    cpy(t,A,(k<<1)+1);
    reverse(t,t+(k<<1)+1);
    cls(t,1,len);
    ntt(t,len,1);
    for(int i=0;i<len;++i)t[i]=(11)t[i]*ib[i]%mod;</pre>
    ntt(t,len,-1);
    cls(t,1,len);
    reverse(t,t+1);
    ntt(t,len,1);
```

```
for(int i=0;i<len;++i)t[i]=(11)t[i]*b[i]%mod;</pre>
    ntt(t,len,-1);
    cls(t,1,len);
    for(int i=0;i<k;++i)A[i]=(A[i]-t[i]+mod)%mod;</pre>
    cls(A,k,len);
  void pow(int*A,int n){
    if(n==1) {cls(A, 0, k+1); A[1]=1; return;}
    pow(A, n>>1);
    int len=1;while(len<=(k<<1))len<<=1;</pre>
    ntt(A,len,1)
    for(int i=0;i<len;++i)A[i]=(11)A[i]*A[i]%mod;</pre>
    ntt(A,len,-1);
    pmod(A);
    if(n&1){
      for(int i=k;i;--i)A[i]=A[i-1];A[0]=0;
      pmod(A);
  }
}
int main(){
  n=rd();k=rd();
  for(int i=1;i<=k;++i)f[i]=(mod+rd())%mod;</pre>
  for(int i=0;i<k;++i)h[i]=(mod+rd())%mod;</pre>
  for(int i=a[k]=b[k]=1;i<=k;++i)</pre>
    a[k-i]=b[k-i]=(mod-f[i])%mod;
  int len=1; while(len<=(k<<1))len<<=1;</pre>
  reverse(a,a+k+1);
  poly::inv(a,ib,len)
  poly::cls(ib,k+1,len);
  poly::ntt(b,len,1);
  poly::ntt(ib,len,1);
  poly::pow(a,n);
  int ans=0:
  for(int i=0;i<k;++i)inc(ans,(11)a[i]*h[i]%mod);</pre>
  printf("%d\n",ans);
  return 0;
}
```

5.13 Chinese Remainder

```
lld crt(lld ans[], lld pri[], int n){
  11d M = 1, ret = 0;
  for(int i=0;i<n;i++) M *= pri[i];</pre>
  for(int i=0;i<n;i++){</pre>
    1ld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
    ret += (ans[i]*(M/pri[i])%M * iv)%M;
    ret %= M;
  return ret;
}
/*
Another:
x = a1 \% m1
x = a2 \% m2
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
0 <= x < lcm(m1, m2)
*/
```

5.14 Berlekamp Massey

```
// x: 1-base, p[]: 0-base
template<size_t N>
vector<llf> BM(llf x[N], size_t n){
    size_t f[N]={0}, t=0;llf d[N];
    vector<llf> p[N];
    for(size_t i=1, b=0;i<=n;++i) {
        for(size_t j=0;j<p[t].size();++j)
            d[i]+=x[i-j-1]*p[t][j];
        if(abs(d[i]-=x[i])<=EPS)continue;
        f[t]=i;if(!t){p[++t].resize(i);continue;}
        vector<llf> cur(i-f[b]-1);
        llf k=-d[i]/d[f[b]];cur.PB(-k);
        for(size_t j=0;j<p[b].size();j++)
            cur.PB(p[b][j]*k);
        if(cur.size()<p[t].size())cur.resize(p[t].size());
        for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];
        if(i-f[b]+p[b].size()>=p[t].size()) b=t;
```

ntt::Transform(fq, s + s);

for (int i = 0; i < s + s; ++i)

```
p[++t]=cur;
                                                                        fv[i] = 1LL * fv[i] * fq[i]%kMod * fq[i]%kMod;
                                                                     ntt::InverseTransform(fv, s + s);
  return p[t];
                                                                     Poly res(s);
                                                                     for (int i = 0; i < s; ++i) {
                                                                       res[i] = kMod - fv[i];
                                                                       if (i < (s >> 1)) {
       NTT
5.15
                                                                          int v = 2 * q[i] % kMod;
                                                                          (res[i] += v) >= kMod ? res[i] -= kMod : 0;
// Remember coefficient are mod P
/* p=a*2^n+1
        2^n
                                        root
   n
                                                                     a = res:
                      65537
   16
        65536
                                                                     if (s >= n) break;
                                        3 */
        1048576
                      7340033
// (must be 2<sup>k</sup>)
                                                                   q.resize(n);
template<LL P, LL root, int MAXN>
                                                                   return q;
struct NTT{
  static LL bigmod(LL a, LL b) {
                                                                 Poly Divide(const Poly &a, const Poly &b) {
    LL res = 1;
                                                                   int n = a.size(), m = b.size(), k = 2;
    for (LL bs = a; b; b >>= 1, bs = (bs * bs) % P)
                                                                   while (k < n - m + 1) k <<= 1;
      if(b&1) res=(res*bs)%P;
                                                                   Poly ra(k), rb(k);
    return res;
                                                                   for (int i = 0; i < min(n, k); ++i) ra[i] = a[n-1-i];
                                                                   for (int i = 0; i < min(m, k); ++i) rb[i] = b[m-1-i];</pre>
  static LL inv(LL a, LL b) {
                                                                   auto rbi = Inverse(rb);
    if(a==1)return 1;
                                                                   auto res = Multiply(rbi, ra);
    return (((LL)(a-inv(b%a,a))*b+1)/a)%b;
                                                                   res.resize(n - m + 1);
                                                                   reverse(res.begin(), res.end());
 LL omega[MAXN+1];
                                                                   return res;
 NTT() {
    omega[0] = 1;
                                                                 Poly Modulo(const Poly &a, const Poly &b) {
    LL r = bigmod(root, (P-1)/MAXN);
                                                                   if (a.size() < b.size()) return a;</pre>
    for (int i=1; i<=MAXN; i++)</pre>
                                                                   auto dv = Multiply(Divide(a, b), b);
      omega[i] = (omega[i-1]*r)%P;
                                                                   assert(dv.size() == a.size());
                                                                   for (int i = 0; i < dv.size(); ++i)
  dv[i] = (a[i] + kMod - dv[i]) % kMod;</pre>
  // n must be 2^k
  void tran(int n, LL a[], bool inv_ntt=false){
                                                                   while (!dv.empty() && dv.back() == 0) dv.pop_back();
    int basic = MAXN / n , theta = basic;
                                                                   return dv;
    for (int m = n; m >= 2; m >>= 1) {
      int mh = m >> 1;
                                                                 Poly Integral(const Poly &f) {
      for (int i = 0; i < mh; i++) {
                                                                   int n = f.size();
        LL w = omega[i*theta%MAXN];
                                                                   VI res(n + 1);
        for (int j = i; j < n; j += m) {</pre>
                                                                   for (int i = 0; i < n; ++i)
           int k = j + mh;
                                                                     res[i+1] = 1LL * f[i] * fpow(i + 1, kMod - 2)%kMod;
           LL x = a[j] - a[k];
                                                                   return res;
           if (x < 0) x += P;
           a[j] += a[k];
                                                                 Poly Evaluate(const Poly &f, const VI &x) {
          if (a[j] > P) a[j] -= P;
                                                                   if (x.empty()) return Poly();
          a[k] = (w * x) % P;
                                                                   int n = x.size();
                                                                   vector<Poly> up(n * 2);
                                                                   for (int i = 0; i < n; ++i) up[i+n] = {kMod-x[i], 1};
for (int i = n - 1; i > 0; --i)
      theta = (theta * 2) % MAXN;
                                                                     up[i] = Multiply(up[i * 2], up[i * 2 + 1]);
    int i = 0;
                                                                   vector<Poly> down(n * 2)
    for (int j = 1; j < n - 1; j++) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1);
                                                                   down[1] = Modulo(f, up[1]);
                                                                   for (int i = 2; i < n * 2; ++i)
      if (j < i) swap(a[i], a[j]);</pre>
                                                                     down[i] = Modulo(down[i >> 1], up[i]);
                                                                   VI y(n);
    if (inv_ntt) {
                                                                   for (int i = 0; i < n; ++i) y[i] = down[i + n][0];
      LL ni = inv(n,P);
                                                                   return y;
      reverse( a+1 , a+n );
for (i = 0; i < n; i++)
                                                                 Poly Interpolate(const VI &x, const VI &y) {
        a[i] = (a[i] * ni) % P;
                                                                   int n = x.size();
                                                                   vector<Poly> up(n * 2);
 }
                                                                   for (int i = 0; i < n; ++i) up[i+n] = {kMod-x[i], 1};
for (int i = n - 1; i > 0; --i)
const LL P=2013265921, root=31;
                                                                     up[i] = Multiply(up[i * 2], up[i * 2 + 1]);
const int MAXN=4194304;
                                                                   VI a = Evaluate(Derivative(up[1]), x);
NTT<P, root, MAXN> ntt;
                                                                   for (int i = 0; i < n; ++i)
a[i] = 1LL * y[i] * fpow(a[i], kMod - 2) % kMod;
                                                                   vector<Poly> down(n * 2);
5.16 Polynomial Operations
                                                                   for (int i = 0; i < n; ++i) down[i + n] = {a[i]};
                                                                   for (int i = n - 1; i > 0; --i) {
  auto lhs = Multiply(down[i * 2], up[i * 2 + 1]);
using VI = vector<int>;
Poly Inverse(Poly f) {
  int n = f.size()
                                                                     auto rhs = Multiply(down[i * 2 + 1], up[i * 2]);
  Poly q(1, fpow(f[0], kMod - 2));
                                                                     assert(lhs.size() == rhs.size());
  for (int s = 2;; s <<= 1) {
  if (f.size() < s) f.resize(s);</pre>
                                                                     down[i].resize(lhs.size());
                                                                     for (int j = 0; j < lhs.size(); ++j)
  down[i][j] = (lhs[j] + rhs[j]) % kMod;</pre>
    Poly fv(f.begin(), f.begin() + s);
    Poly fq(q.begin(), q.end());
fv.resize(s + s); fq.resize(s + s);
                                                                   return down[1];
    ntt::Transform(fv, s + s);
```

Poly Log(Poly f) {

int n = f.size();

```
if( x[ i ] >= MOD ) x[ i ] -= MOD;
if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
  if (n == 1) return {0};
  auto d = Derivative(f);
  f.resize(n - 1);
  d = Multiply(d, Inverse(f));
                                                                if( inv )
  d.resize(n - 1)
                                                                  for( int i = 0 ; i < N ; i++ ) {</pre>
  return Integral(d);
                                                                    x[ i ] *= inv( N, MOD );
                                                                    x[ i ] %= MOD;
Poly Exp(Poly f) {
  int n = f.size()
  Poly q(1, 1); f[0] += 1;
                                                              }
  for (int s = 1; s < n; s <<= 1) {
    if (f.size() < s + s) f.resize(s + s);</pre>
                                                              5.18 DiscreteLog
    Poly g(f.begin(), f.begin() + s + s);
    Poly h(q.begin(), q.end());
h.resize(s + s); h = Log(h);
                                                              // Baby-step Giant-step Algorithm
                                                              11d BSGS(11d P, 11d B, 11d N) {
    for (int i = 0; i < s + s; ++i)
                                                                // find B^L = N \mod P
      g[i] = (g[i] + kMod - h[i]) % kMod;
                                                                unordered_map<lld, int> R;
    g = Multiply(g, q);
                                                                11d sq = (11d)sqrt(P);
    g.resize(s + s); q = g;
                                                                11d t = 1;
                                                                for (int i = 0; i < sq; i++) {
  assert(q.size() >= n);
                                                                  if (t == N) return i
  q.resize(n);
                                                                  if (!R.count(t)) R[t] = i;
  return q;
                                                                  t = (t * B) % P;
Poly SquareRootImpl(Poly f) {
                                                                11d f = inverse(t, P);
  if (f.empty()) return {0};
                                                                for(int i=0;i<=sq+1;i++) {</pre>
  int z = QuadraticResidue(f[0], kMod), n = f.size();
                                                                  if (R.count(N))
  constexpr int kInv2 = (kMod + 1) >> 1;
                                                                    return i * sq + R[N];
  if (z == -1) return {-1};
                                                                  N = (N * f) % P;
  VI q(1, z);
  for (int s = 1; s < n; s <<= 1) {
                                                                return -1:
    if (f.size() < s + s) f.resize(s + s);</pre>
    VI fq(q.begin(), q.end());
    fq.resize(s + s)
    VI f2 = Multiply(fq, fq);
                                                              5.19 Quadratic residue
    f2.resize(s + s);
    for (int i = 0; i < s + s; ++i)
                                                              struct Status{
      f2[i] = (f2[i] + kMod - f[i]) % kMod;
                                                                11 x,y;
    f2 = Multiply(f2, Inverse(fq));
    f2.resize(s + s);
                                                              11 w;
    for (int i = 0; i < s + s; ++i)
                                                              Status mult(const Status& a,const Status& b,ll mod){
      fq[i] = (fq[i]+kMod - 1LL*f2[i]*kInv2%kMod)%kMod;
                                                                Status res;
                                                                res.x=(a.x*b.x+a.y*b.y%mod*w)%mod;
                                                                res.y=(a.x*b.y+a.y*b.x)%mod;
  q.resize(n);
                                                                return res;
  return q;
                                                              inline Status qpow(Status _base, 11 _pow, 11 _mod) {
Poly SquareRoot(Poly f) {
                                                                Status res = \{1, 0\};
  int n = f.size(), m = 0;
while (m < n && f[m] == 0) m++;</pre>
                                                                while(_pow>0){
                                                                  if(_pow&1) res=mult(res,_base,_mod);
  if (m == n) return VI(n);
                                                                  _base=mult(_base,_base,_mod);
  if (m & 1) return {-1};
                                                                  _pow>>=1;
  auto s = SquareRootImpl(VI(f.begin() + m, f.end()));
  if (s[0] == -1) return {-1};
                                                                return res;
  VI res(n);
  for (int i = 0; i < s.size(); ++i) res[i + m/2]=s[i];</pre>
                                                              inline 11 check(11 x,11 p){
  return res;
                                                                return qpow_mod(x,(p-1)>>1,p);
                                                              inline 11 get_root(11 n,11 p){
                                                                if(p==2) return 1;
5.17 FWT
                                                                if(check(n,p)==p-1) return -1;
                                                                11 a;
/* xor convolution:
                                                                while(true){
                                                                  a=rand()%p;
```

```
* x = (x0,x1) , y = (y0,y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
* x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x\theta+x1)(y\theta+y1)), (x\theta-x1)(y\theta-y1))
* z = (1/2) * z''
 * or convolution:
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
* and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
const LL MOD = 1e9+7;
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
  for( int d = 1 ; d < N ; d <<= 1 ) {
    int d2 = d << 1;
    for( int s = 0 ; s < N ; s += d2 )
      for( int i = s , j = s+d ; i < s+d ; i++, j++ ){</pre>
        LL ta = x[i], tb = x[j];
        x[ i ] = ta+tb;
        x[ j ] = ta-tb;
```

5.20 De-Bruijn

Status res = {a, 1} res=qpow(res,(p+1)>>1,p);

return res.x;

w=((a*a-n)%p+p)%p;

if(check(w,p)==p-1) break;

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
  if (t > n) {
    if (n % p == 0)
      for (int i = 1; i <= p; ++i)
      res[sz++] = aux[i];
  } else {</pre>
```

```
aux[t] = aux[t - p];
    db(t + 1, p, n, k)
    for (int i = aux[t - p] + 1; i < k; ++i) {
     aux[t] = i;
     db(t + 1, t, n, k);
 }
int de_bruijn(int k, int n) {
 // return cyclic string of len k^n s.t. every string
  // of len n using k char appears as a substring.
 if (k == 1) {
    res[0] = 0;
    return 1;
 for (int i = 0; i < k * n; i++) aux[i] = 0;
 sz = 0:
 db(1, 1, n, k);
  return sz;
```

5.21 Simplex Construction

Standard form: maximize $\sum_{1\leq i\leq n}c_ix_i$ such that for all $1\leq j\leq m$, $\sum_{1\leq i\leq n}A_{ji}x_i\leq b_j$ and $x_i\geq 0$ for all $1\leq i\leq n$.

- 1. In case of minimization, let $c_i^\prime = -c_i$
- 2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \rightarrow \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- $3. \sum_{1 \leq i \leq n} A_{ji} x_i = b_j$
 - $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$
 - $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

5.22 Simplex

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return vector<double>(n, -inf) if the solution doesn
     't exist
// return vector<double>(n, +inf) if the solution is
    unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d:
vector<int> p, q;
void pivot(int r, int s) {
  double inv = 1.0 / d[r][s];
  for (int i = 0; i < m + 2; ++i) {</pre>
    for (int j = 0; j < n + 2; ++j) {
  if (i != r && j != s)</pre>
        d[i][j] -= d[r][j] * d[i][s] * inv;
    }
  for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
  for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;
d[r][s] = inv;</pre>
  swap(p[r], q[s]);
bool phase(int z) {
  int x = m + z;
  while (true) {
    int s = -1;
    for (int i = 0; i <= n; ++i) {</pre>
      if (!z && q[i] == -1) continue;
      if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
    if (d[x][s] > -eps) return true;
    for (int i = 0; i < m; ++i) {
      if (d[i][s] < eps) continue;</pre>
        d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
    if (r == -1) return false;
    pivot(r, s);
```

```
}
VD solve(const VVD &a, const VD &b, const VD &c) {
  m = b.size(), n = c.size();
  d = VVD(m + 2, VD(n + 2));
  for (int i = 0; i < m; ++i) {
     for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
  p.resize(m), q.resize(n + 1);
  for (int i = 0; i < m; ++i)
  p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
  q[n] = -1, d[m + 1][n] = 1;
   int r = 0;
  for (int i = 1; i < m; ++i)
     if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
  if (d[r][n + 1] < -eps) {</pre>
     pivot(r, n);
     if (!phase(1) || d[m + 1][n + 1] < -eps)
     return VD(n, -inf);
for (int i = 0; i < m; ++i) if (p[i] == -1) {</pre>
       int s = min_element(d[i].begin(), d[i].end() - 1)
                - d[i].begin();
       pivot(i, s);
  if (!phase(0)) return VD(n, inf);
  VD x(n);
  for (int i = 0; i < m; ++i)</pre>
     if (p[i] < n) x[p[i]] = d[i][n + 1];
  return x;
}}
```

6 Geometry

6.1 Point Class

```
template<typename T>
struct Point{
  typedef long double 11f:
  static constexpr llf EPS = 1e-8;
  T x, y;
  Point(T _=0, T _{_}=0): x(_), y(_{_})\{\}
  template<typename T2>
    Point(const Point<T2>& a): x(a.x), y(a.y){}
  inline llf theta() const {
    return atan2((11f)y, (11f)x);}
  inline llf dis() const
    return hypot((llf)x, (llf)y);}
  inline 11f dis(const Point& o) const {
    return hypot((llf)(x-o.x), (llf)(y-o.y));}
  Point operator-(const Point& o) const {
    return Point(x-o.x, y-o.y);}
  Point operator+(const Point& o) const {
  return Point(x+o.x, y+o.y);}
Point operator*(const T& k) const {
    return Point(x*k, y*k);}
  Point operator/(const T& k) const {
    return Point(x/k, y/k);}
  Point operator-() const {return Point(-x, -y);}
  Point rot90() const {return Point(-y, x);}
  template<typename T2>
  bool in(const Circle<T2>& a) const {
    /* Add struct Circle at top */
    return a.o.dis(*this)+EPS <= a.r;</pre>
  bool equal(const Point& o, true_type) const {
    return fabs(x-o.x) < EPS and fabs(y-o.y) < EPS; }</pre>
  bool equal(const Point& o, false_type) const {
    return tie(x, y) == tie(o.x, o.y); }
  bool operator==(const Point& o) const {
    return equal(o, is_floating_point<T>()); }
  bool operator!=(const Point& o) const {
    return !(*this == o); }
  bool operator<(const Point& o) const {</pre>
    return theta() < o.theta();</pre>
    // sort like what pairs did
    // if (is_floating_point<T>())
         return fabs(x-o.x)<EPS?y<o.y:x<o.x;
    // else return tie(x, y) < tie(o.x, o.y);</pre>
  friend inline T cross(const Point&a,const Point&b){
```

```
return a.x*b.y - b.x*a.y; }
friend inline T dot(const Point& a, const Point &b){
  return a.x*b.x + a.y*b.y; }
friend ostream&operator<<(ostream&ss,const Point&o){
    ss<<"("<<0.x<<", "<<0.y<<")"; return ss; }
};</pre>
```

6.2 Circle Class

```
template<typename T>
struct Circle{
    static constexpr llf EPS = 1e-8;
    Point<T> o; T r;
    vector<Point<llf>> operator&(const Circle& aa)const{
        //https://www.cnblogs.com/wangzming/p/8338142.html
        llf d=o.dis(aa.o);
        if(d>r+aa.r+EPS || d<fabs(r-aa.r)-EPS) return {};
        llf dt = (r*r - aa.r*aa.r)/d, d1 = (d+dt)/2;
        Point<llf>> dir = (aa.o-o); dir /= d;
        Point<llf>> pcrs = dir*d1 + o;
        dt=sqrt(max(0.0L, r*r - d1*d1)), dir=dir.rot90();
        return {pcrs + dir*dt, pcrs - dir*dt};
    }
};
```

6.3 Segment Class

```
const long double EPS = 1e-8;
template<typename T>
struct Segment{
  // p1.x < p2.x
 Line<T> base;
 Point<T> p1, p2;
 Segment(): base(Line<T>()), p1(Point<T>()), p2(Point<</pre>
      T>()){
    assert(on_line(p1, base) and on_line(p2, base));
 Segment(Line<T> _, Point<T> __, Point<T> ___): base(_
      ), p1(__), p2(___){
    assert(on_line(p1, base) and on_line(p2, base));
 template<typename T2>
    Segment(const Segment<T2>& _): base(_.base), p1(_.
        p1), p2(_.p2) {}
  typedef Point<long double> Pt;
  friend bool on_segment(const Point<T>& p, const
      Segment& 1){
    if(on_line(p, l.base))
      return (1.p1.x-p.x)*(p.x-1.p2.x)>=0 and (1.p1.y-p
         .y)*(p.y-1.p2.y)>=0;
    return false;
  friend bool have_inter(const Segment& a, const
      Seament& b){
    if(is_parallel(a.base, b.base)){
      return on_segment(a.p1, b) or on_segment(a.p2, b)
          or on_segment(b.p1, a) or on_segment(b.p2, a
          );
   Pt inter = get_inter(a.base, b.base);
    return on_segment(inter, a) and on_segment(inter, b
  friend inline Pt get_inter(const Segment& a, const
     Segment& b){
    if(!have_inter(a, b)){
      return NOT_EXIST;
    }else if(is_parallel(a.base, b.base)){
     if(a.p1 == b.p1){
        if(on_segment(a.p2, b) or on_segment(b.p2, a))
            return INF_P;
        else return a.p1;
      else if(a.p1 == b.p2){
        if(on_segment(a.p2, b) or on_segment(b.p1, a))
            return INF_P;
        else return a.p1;
      }else if(a.p2 == b.p1){
        if(on_segment(a.p1, b) or on_segment(b.p2, a))
            return INF_P;
        else return a.p2;
      else if(a.p2 == b.p2){
```

6.4 Line Class

```
const Point<long double> INF_P(-1e20, 1e20);
const Point<long double> NOT_EXIST(1e20, 1e-20);
template<typename T>
struct Line{
 static constexpr long double EPS = 1e-8;
  // ax+by+c = 0
 T a, b, c;
Line(T _=0, T __=1, T ___=0): a(_), b(__), c(___){
    assert(fabs(a)>EPS or fabs(b)>EPS);}
  template<typename T2>
    Line(const Line<T2>& x): a(x.a), b(x.b), c(x.c){}
  typedef Point<long double> Pt;
  bool equal(const Line& o, true_type) const {
    return fabs(a-o.a)<EPS &&
    fabs(b-o.b)<EPS && fabs(c-o.b)<EPS;}</pre>
  bool equal(const Line& o, false_type) const {
    return a==o.a and b==o.b and c==o.c;}
  bool operator==(const Line& o) const {
    return equal(o, is_floating_point<T>());}
  bool operator!=(const Line& o) const {
    return !(*this == o);}
  friend inline bool on_line__(const Point<T>& p, const
    Line& 1, true_type) {
return fabs(1.a*p.x + 1.b*p.y + 1.c) < EPS;
  friend inline bool on_line__(const Point<T>& p, const
       Line& 1, false_type){
    return 1.a*p.x + 1.b*p.y + 1.c == 0;
  friend inline bool on_line(const Point<T>&p, const
      Line& 1){
    return on_line__(p, l, is_floating_point<T>());
  friend inline bool is_parallel__(const Line& x, const
       Line& y, true_type){
    return fabs(x.a*y.b - x.b*y.a) < EPS;</pre>
  friend inline bool is_parallel__(const Line& x, const
    Line& y, false_type){
return x.a*y.b == x.b*y.a;
  friend inline bool is_parallel(const Line& x, const
      Line& y){
    return is_parallel__(x, y, is_floating_point<T>());
  friend inline Pt get_inter(const Line& x, const Line&
    typedef long double llf;
    if(x==y) return INF_P;
    if(is_parallel(x, y)) return NOT_EXIST;
    llf delta = x.a*y.b - x.b*y.a;
    llf delta_x = x.b*y.c - x.c*y.b;
    11f delta_y = x.c*y.a - x.a*y.c;
    return Pt(delta_x / delta, delta_y / delta);
  friend ostream&operator<<(ostream&ss, const Line&o){</pre>
    ss<<o.a<<"x+"<<o.b<<"y+"<<o.c<<"=0";
    return ss;
```

```
};
template<typename T>
inline Line<T> get_line(const Point<T>& a, const Point<
    T>& b){
    return Line<T>(a.y-b.y, b.x-a.x, (b.y-a.y)*a.x-(b.x-a.x)*a.y);
}
```

6.5 Triangle Circumcentre

6.6 2D Convex Hull

```
template<typename T>
class ConvexHull_2D{
private:
 typedef Point<T> PT;
  vector<PT> d;
 struct myhash{
    uint64_t operator()(const PT& a) const {
      uint64_t xx=0, yy=0;
      memcpy(&xx, &a.x, sizeof(a.x));
      memcpy(&yy, &a.y, sizeof(a.y));
      uint64_t ret = xx*17+yy*31;
      ret = (ret ^ (ret >> 16))*0x9E3779B1;
      ret = (ret ^ (ret >> 13))*0xC2B2AE35;
      ret = ret ^ xx;
      return (ret ^ (ret << 3)) * yy;
 unordered_set<PT, myhash> in_hull;
public:
 void init(){in_hull.clear();d.clear();}
  void insert(const PT& x){d.PB(x);}
 void solve(){
    sort(ALL(d), [](const PT& a, const PT& b){
      return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
    vector<PT> s(SZ(d) << 1); int o = 0;
    for(auto p: d) {
      while(o>=2 \&\& cross(p-s[o-2],s[o-1]-s[o-2])<=0)
        0--
      s[o++] = p;
    for(int i=SZ(d)-2, t = o+1;i>=0;i--){
      while(o = t\&cross(d[i] - s[o-2], s[o-1] - s[o-2]) <= 0)
        0--
      s[o++] = d[i];
    s.resize(o-1); swap(s, d);
    for(auto i: s) in_hull.insert(i);
 vector<PT> get(){return d;}
 bool in_it(const PT& x){
    return in_hull.find(x)!=in_hull.end();}
```

6.7 2D Farthest Pair

```
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {
   while(abs(cross(stk[i+1]-stk[i],
        stk[(pos+1)%n]-stk[i])) >
        abs(cross(stk[i+1]-stk[i],
        stk[pos]-stk[i]))) pos = (pos+1)%n;
ans = max({ans, dis(stk[i], stk[pos]),
```

```
6.8 2D Closest Pair
```

dis(stk[i+1], stk[pos])});

```
struct cmp_y {
  bool operator()(const P& p, const P& q) const {
    return p.y < q.y;</pre>
  }
};
multiset<P, cmp_y> s;
void solve(P a[], int n) {
  sort(a, a + n, [](const P& p, const P& q) {
    return tie(p.x, p.y) < tie(q.x, q.y);
  11f d = INF; int pt = 0;
  for (int i = 0; i < n; ++i) {
    while (pt < i \text{ and } a[i].x - a[pt].x >= d)
      s.erase(s.find(a[pt++]));
    auto it = s.lower_bound(P(a[i].x, a[i].y - d));
    while (it != s.end() and it->y - a[i].y < d)</pre>
      d = min(d, dis(*(it++), a[i]));
    s.insert(a[i]);
}
```

6.9 Simulated Annealing

```
11f anneal() {
  mt19937 rnd_engine( seed );
  uniform_real_distribution< llf > rnd( 0, 1 );
  const llf dT = 0.001;
   // Argument p
  llf S_cur = calc( p ), S_best = S_cur;
for ( llf T = 2000 ; T > EPS ; T -= dT ) {
     // Modify p to p_prime
     const 11f S_prime = calc( p_prime );
     const llf delta_c = S_prime - S_cur;
     11f prob = min( ( llf ) 1, exp( -delta_c / T ) );
     if ( rnd( rnd_engine ) <= prob )</pre>
       S_cur = S_prime, p = p_prime;
     if ( S_prime < S_best )</pre>
       S_best = S_prime, p_best = p_prime;
  return S_best;
}
```

6.10 Half Plane Intersection

```
inline int dcmp ( double x ) {
  if( fabs( x ) < eps ) return 0;</pre>
  return x > 0 ? 1 : -1;
struct Line {
  Point st, ed;
  double ang;
  Line(Point _s=Point(), Point _e=Point()):
   st(_s),ed(_e),ang(atan2(_e.y-_s.y,_e.x-_s.x)){}
  inline bool operator< ( const Line& rhs ) const {</pre>
    if(dcmp(ang - rhs.ang) != 0) return ang < rhs.ang;</pre>
    return dcmp( cross( st, ed, rhs.st ) ) < 0;</pre>
};
// cross(pt, line.ed-line.st)>=0 <-> pt in half plane
vector< Line > lns;
deque< Line > que;
deque< Point > pt;
double HPI() {
  sort( lns.begin(), lns.end() );
  que.clear(); pt.clear();
  que.push_back( lns[ 0 ] );
  for ( int i = 1 ; i < (int)lns.size() ; i ++ ) {</pre>
    if(!dcmp(lns[i].ang - lns[i-1].ang)) continue;
    while ( pt.size() > 0 &&
     dcmp(cross(lns[i].st,lns[i].ed,pt.back()))<0){</pre>
      pt.pop_back();que.pop_back();
    while ( pt.size() > 0 &&
     dcmp(cross(lns[i].st,lns[i].ed,pt.front()))<0){</pre>
      pt.pop_front(); que.pop_front();
```

```
pt.push_back(get_point( que.back(), lns[ i ] ));
  que.push_back( lns[ i ] );
while ( pt.size() > 0 &&
 dcmp(cross(que[0].st, que[0].ed, pt.back()))<0){</pre>
  que.pop_back();
  pt.pop_back();
while ( pt.size() > 0 &&
 dcmp(cross(que.back().st,que.back().ed,pt[0]))<0){</pre>
  que.pop_front();
  pt.pop_front();
pt.push_back(get_point(que.front(), que.back()));
vector< Point > conv;
for ( int i = 0 ; i < (int)pt.size() ; i ++ )</pre>
  conv.push_back( pt[ i ] );
double ret = 0:
for ( int i = 1 ; i + 1 < (int)conv.size() ; i ++ )</pre>
  ret += abs(cross(conv[0], conv[i], conv[i + 1]));
return ret / 2.0;
```

6.11 Ternary Search on Integer

```
int TernarySearch(int 1, int r) {
   // max value @ (1, r]
   while (r - 1 > 1) {
     int m = (1 + r) >> 1;
     if (f(m) > f(m + 1)) r = m;
     else 1 = m;
   }
   return 1+1;
}
```

6.12 Minimum Covering Circle

```
template<typename T>
Circle<llf> MinCircleCover(const vector<PT>& pts){
  random_shuffle(ALL(pts));
  Circle<llf> c = \{pts[0], 0\};
  for(int i=0;i<SZ(pts);i++){</pre>
    if(pts[i].in(c)) continue;
    c = {pts[i], 0};
    for(int j=0;j<i;j++){</pre>
      if(pts[j].in(c)) continue;
      c.o = (pts[i] + pts[j]) / 2;
      c.r = pts[i].dis(c.o);
      for(int k=0;k<j;k++){</pre>
        if(pts[k].in(c)) continue;
        c = get_circum(pts[i], pts[j], pts[k]);
   }
  return c;
```

6.13 KDTree (Nearest Point)

```
const int MXN = 100005;
struct KDTree {
  struct Node {
    int x,y,x1,y1,x2,y2;
    int id,f;
Node *L, *R;
  } tree[MXN], *root;
  int n:
 LL dis2(int x1, int y1, int x2, int y2) {
  LL dx = x1-x2, dy = y1-y2;
    return dx*dx+dy*dy;
  static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
  static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
  void init(vector<pair<int,int>> ip) {
    n = ip.size();
    for (int i=0; i<n; i++) {
      tree[i].id = i;
      tree[i].x = ip[i].first;
      tree[i].y = ip[i].second;
```

```
root = build_tree(0, n-1, 0);
  Node* build_tree(int L, int R, int d) {
    if (L>R) return nullptr
    int M = (L+R)/2; tree[M].f = d%2;
    nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
    tree[M].x1 = tree[M].x2 = tree[M].x;
    tree[M].y1 = tree[M].y2 = tree[M].y;
    tree[M].L = build_tree(L, M-1, d+1);
    if (tree[M].L) {
      tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
    tree[M].R = build_tree(M+1, R, d+1);
    if (tree[M].R) {
      tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
       tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
    return tree+M;
  int touch(Node* r, int x, int y, LL d2){
    LL dis = sqrt(d2)+1;
    if (x<r->x1-dis || x>r->x2+dis ||
        y<r->y1-dis || y>r->y2+dis)
       return 0;
    return 1:
  void nearest(Node* r,int x,int y,int &mID,LL &md2) {
    if (!r || !touch(r, x, y, md2)) return;
    LL d2 = dis2(r->x, r->y, x, y);
    if (d2 < md2 \mid \mid (d2 == md2 && mID < r->id)) {
      mID = r->id;
      md2 = d2;
    // search order depends on split dim
    if ((r->f == 0 \&\& x < r->x) ||
         (r->f == 1 && y < r->y))
       nearest(r->L, x, y, mID, md2);
      nearest(r->R, x, y, mID, md2);
    } else {
      int query(int x, int y) {
    int id = 1029384756;
    LL d2 = 102938475612345678LL;
    nearest(root, x, y, id, d2);
    return id;
} tree;
```

7 Stringology

7.1 Hash

```
class Hash{
private:
  const int p = 127, q = 1051762951;
int sz, prefix[N], power[N];
   int add(int x, int y){return x+y>=q?x+y-q:x+y;}
   int sub(int x, int y){return x-y<0?x-y+q:x-y;}</pre>
   int mul(int x, int y){return 1LL*x*y%q;}
public:
   void init(const string &x){
     sz = x.size();prefix[0]=0;power[0]=1;
     for(int i=1;i<=sz;i++)</pre>
       prefix[i]=add(mul(prefix[i-1], p), x[i-1]);
     for(int i=1;i<=sz;i++)power[i]=mul(power[i-1], p);</pre>
   int query(int 1, int r){
     // 1-base (1, r]
     return sub(prefix[r], mul(prefix[l], power[r-l]));
};
```

7.2 Suffix Array

```
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
int x[maxn], p[maxn], q[maxn * 2];
// sa[i]: sa[i]-th suffix is the \
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix \
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *sa, int *c, int n, int z) {
   memset(sa, 0, sizeof(int) * n);
   memcpy(x, c, sizeof(int) * z);
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
   memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
      if (sa[i] && !t[sa[i] - 1])
        sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
   memcpy(x, c, sizeof(int) * z);
   for (int i = n - 1; i >= 0; --i)
     if (sa[i] && t[sa[i] - 1])
        sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q,
bool *t, int *c, int n, int z) {
  bool uniq = t[n - 1] = true;
   int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
   memset(c, 0, sizeof(int) * z);
for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
   if (uniq) {
     for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
     return;
   for (int i = n - 2; i \ge 0; --i)
     t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
   pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i)</pre>
     if (t[i] && !t[i - 1])
        sa[--x[s[i]]] = p[q[i] = nn++] = i;
   induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) {
     if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
     bool neq = last < 0 || '</pre>
      memcmp(s + sa[i], s + last,
(p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
     ns[q[last = sa[i]]] = nmxz += neq;
   sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
   pre(sa, c, n, z);
for (int i = nn - 1; i >= 0; --i)
     sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
   induce(sa, c, s, t, n, z);
void build(const string &s) {
   for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];
   _s[(int)s.size()] = 0; // s shouldn't contain 0
   sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];
for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
   int ind = 0; hi[0] = 0;
   for (int i = 0; i < (int)s.size(); ++i) {</pre>
     if (!rev[i]) {
        ind = 0:
        continue;
     while (i + ind < (int)s.size() && \</pre>
       s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
     hi[rev[i]] = ind ? ind-- : 0;
}}
```

7.3 Aho-Corasick Algorithm

```
class AhoCorasick{
  private:
    static constexpr int Z = 26;
    struct node{
      node *nxt[ Z ], *fail;
```

```
vector< int > data:
      node(): fail( nullptr ) {
        memset( nxt, 0, sizeof( nxt ) );
        data.clear();
    } *rt;
    inline int Idx( char c ) { return c - 'a'; }
  public:
    void init() { rt = new node(); }
    void add( const string& s, int d ) {
      node* cur = rt;
      for ( auto c : s ) {
        if ( not cur->nxt[ Idx( c ) ] )
          cur->nxt[ Idx( c ) ] = new node();
        cur = cur->nxt[ Idx( c ) ];
      cur->data.push back( d ):
    void compile() {
      vector< node* > bfs;
      size_t ptr = 0;
      for ( int i = 0 ; i < Z ; ++ i ) {
        if ( not rt->nxt[_i ] ) {
          // uncomment 2 lines to make it DFA
          // rt->nxt[i] = rt;
          continue;
        rt->nxt[ i ]->fail = rt;
        bfs.push_back( rt->nxt[ i ] );
      while ( ptr < bfs.size() ) {
  node* u = bfs[ ptr ++ ];</pre>
        for ( int i = 0 ; i < Z ; ++ i ) {
          if ( not u->nxt[ i ] ) {
            // u->nxt[i] = u->fail->nxt[i];
            continue;
          node* u_f = u->fail;
          while ( u_f ) {
            if ( not u_f->nxt[ i ] ) {
              u_f = u_f->fail; continue;
            u->nxt[ i ]->fail = u_f->nxt[ i ];
            break:
           if ( not u_f ) u->nxt[ i ]->fail = rt;
          bfs.push_back( u->nxt[ i ] );
        }
    void match( const string& s, vector< int >& ret ) {
      node* u = rt;
      for ( auto c : s ) {
        while ( u != rt and not u->nxt[ Idx( c ) ] )
          u = u->fail
        u = u->nxt[ Idx( c ) ];
        if ( not u ) u = rt;
        node* tmp = u;
        while ( tmp != rt ) {
          for ( auto d : tmp->data )
            ret.push_back( d );
          tmp = tmp->fail;
      }
} ac;
```

7.4 Suffix Automaton

```
struct Node{
  Node *green, *edge[26];
  int max_len;
  Node(const int _max_len)
    : green(NULL), max_len(_max_len){
     memset(edge,0,sizeof(edge));
  }
} *ROOT, *LAST;
void Extend(const int c) {
  Node *cursor = LAST;
  LAST = new Node((LAST->max_len) + 1);
  for(;cursor&&!cursor->edge[c]; cursor=cursor->green)
```

int i, j, left, right;

left=right=0; z[0]=len;
for(i=1;i<len;i++) {</pre>

```
cursor->edge[c] = LAST;
                                                                  j=max(min(z[i-left],right-i),0);
  if (!cursor)
                                                                  for(;i+j<len&&s[i+j]==s[j];j++);
                                                                  z[i]=j;
    LAST->green = ROOT;
                                                                  if(i+z[i]>right) {
  else {
    Node *potential_green = cursor->edge[c];
                                                                    right=i+z[i];
    if((potential_green->max_len)==(cursor->max_len+1))
                                                                    left=i;
      LAST->green = potential_green;
    else {
//assert(potential_green->max_len>(cursor->max_len+1));
      Node *wish = new Node((cursor->max_len) + 1);
      for(;cursor && cursor->edge[c]==potential_green;
           cursor = cursor->green)
                                                             7.7 Manacher
        cursor->edge[c] = wish;
      for (int i = 0; i < 26; i++)
  wish->edge[i] = potential_green->edge[i];
                                                             int z[maxn];
                                                             int manacher(const string& s) {
  string t = ".";
      wish->green = potential_green->green;
      potential_green->green = wish;
                                                                for(char c:s)) t += c, t += '.';
      LAST->green = wish;
                                                                int 1 = 0, r = 0, ans = 0;
                                                                for (int i = 1; i < t.length(); ++i) {</pre>
  }
                                                                  z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
                                                                  while (i - z[i] \ge 0 \&\& i + z[i] < t.length()) {
char S[10000001], A[10000001];
                                                                    if(t[i - z[i]] == t[i + z[i]]) ++z[i];
int N:
int main(){
                                                                    else break:
  scanf("%d%s", &N, S);
  ROOT = LAST = new Node(0)
                                                                  if (i + z[i] > r) r = i + z[i], l = i;
  for (int i = 0; S[i]; i++)
    Extend(S[i] - 'a');
                                                                for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
  while (N--){
                                                                return ans;
    scanf("%s", A);
    Node *cursor = ROOT;
    bool ans = true;
    for (int i = 0; A[i]; i++){
                                                                     Lexico Smallest Rotation
      cursor = cursor->edge[A[i] - 'a'];
      if (!cursor) {
        ans = false;
                                                             string mcp(string s){
                                                                int n = s.length();
        break;
                                                                s += s;
    }
                                                                int i=0, j=1;
                                                                while (i<n && j<n){</pre>
    puts(ans ? "Yes" : "No");
                                                                  int k = 0;
                                                                  while (k < n \&\& s[i+k] == s[j+k]) k++;
  return 0;
                                                                  if (s[i+k] \le s[j+k]) j += k+1;
                                                                  else i += k+1;
                                                                  if (i == j) j++;
7.5 KMP
                                                                int ans = i < n ? i : j;</pre>
vector<int> kmp(const string &s) {
                                                                return s.substr(ans, n);
  vector<int> f(s.size(), 0);
  /* f[i] = length of the longest prefix
     (excluding s[0:i]) such that it coincides
     with the suffix of s[0:i] of the same length */
                                                             7.9
                                                                     BWT
   * i + 1 - f[i] is the length of the
     smallest recurring period of s[0:i] */
  int k = 0;
                                                             struct BurrowsWheeler{
  for (int i = 1; i < (int)s.size(); ++i) {</pre>
                                                             #define SIGMA 26
    while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
                                                             #define BASE 'a'
    if (s[i] == s[k]) ++k;
                                                                vector<int> v[ SIGMA ];
    f[i] = k;
                                                                void BWT(char* ori, char* res){
  }
                                                                  // make ori -> ori + ori
  return f;
                                                                  // then build suffix array
vector<int> search(const string &s, const string &t) {
                                                                void iBWT(char* ori, char* res){
  // return 0-indexed occurrence of t in s
                                                                  for( int i = 0 ; i < SIGMA ; i ++ )</pre>
  vector<int> f = kmp(t), r;
for (int i = 0, k = 0; i < (int)s.size(); ++i)</pre>
                                                                    v[ i ].clear();
                                                                  int len = strlen( ori );
    while(k > 0 && (k==(int)t.size() || s[i]!=t[k]))
                                                                  for( int i = 0 ; i < len ; i ++</pre>
      k = f[k - 1]
                                                                    v[ ori[i] - BASE ].push_back( i );
    if (s[i] == t[k]) ++k;
                                                                  vector<int> a;
    if (k == (int)t.size()) r.push_back(i-t.size()+1);
                                                                  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )
for( auto j : v[ i ] ){</pre>
  return res;
                                                                      a.push_back( j );
                                                                      ori[ ptr ++ ] = BASE + i;
                                                                  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){</pre>
7.6 Z value
                                                                    res[ i ] = ori[ a[ ptr ] ];
char s[MAXN];
                                                                    ptr = a[ ptr ];
int len,z[MAXN];
void Z_value() {
                                                                  res[ len ] = 0;
```

} bwt;

7.10 Palindromic Tree

```
struct palindromic_tree{
  struct node{
    int next[26],f,len;
    int cnt,num,st,ed;
    node(int 1=0):f(0),len(1),cnt(0),num(0) {
      memset(next, 0, sizeof(next)); }
  vector<node> st;
  vector<char> s;
  int last,n;
  void init(){
    st.clear();s.clear();last=1; n=0;
    st.push_back(0);st.push_back(-1);
    st[0].f=1;s.push_back(-1); }
  int getFail(int x){
    while(s[n-st[x].len-1]!=s[n])x=st[x].f;
    return x;}
  void add(int c){
    s.push_back(c-='a'); ++n;
    int cur=getFail(last);
    if(!st[cur].next[c]){
      int now=st.size();
      st.push_back(st[cur].len+2);
      st[now].f=st[getFail(st[cur].f)].next[c];
      st[cur].next[c]=now;
      st[now].num=st[st[now].f].num+1;
    last=st[cur].next[c];
    ++st[last].cnt;}
  int size(){ return st.size()-2;}
} pt;
int main() {
  string s; cin >> s; pt.init();
  for (int i=0; i<SZ(s); i++) {</pre>
    int prvsz = pt.size(); pt.add(s[i]);
    if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
      // pal @ [1,r]: s.substr(1, r-1+1)
    }
  return 0:
```

8 Misc

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G , where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G .

- The number of undirected spanning in G is $|\mathsf{det}(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|{\rm det}(\tilde{L}_{rr})|$.

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\ldots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all $1 \leq k \leq n$.

8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W\subseteq X, |W|\leq |N_G(W)|$

8.1.7 Euler's planar graph formula

```
V - E + F = C + 1, E \le 3V - 6(?)
```

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.2 MaximumEmptyRect

```
int max_empty_rect(int n, int m, bool blocked[N][N]) {
  static int mxu[2][N], me=0, he=1, ans=0;
  for (int i=0;i<m;i++) mxu[he][i]=0;
for (int i=0;i<n;i++) {</pre>
    stack<PII, vector<PII>> stk;
    for (int j=0;j<m;++j) {
  if (blocked[i][j]) mxu[me][j]=0;</pre>
      else mxu[me][j]=mxu[he][j]+1;
      int la = j;
      while (!stk.empty()&&stk.top().FF>mxu[me][j]) {
         int x1 = i - stk.top().FF, x2 = i;
         int y1 = stk.top().SS, y2 = j;
         la = stk.top().SS; stk.pop();
         ans=\max(ans,(x2-x1)*(y2-y1));
      if (stk.empty()||stk.top().FF<mxu[me][j])</pre>
         stk.push({mxu[me][j],la});
    while (!stk.empty()) {
      int x1 = i - stk.top().FF, x2 = i;
       int y1 = stk.top().SS-1, y2 = m-1;
      stk.pop(); ans=max(ans,(x2-x1)*(y2-y1));
    swap(me, he);
  }
  return ans;
}
```

8.3 DP-opt Condition

8.3.1 totally monotone (concave/convex)

```
\begin{array}{c} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

8.3.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', \ B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', \ B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

8.4 Convex 1D/1D DP

```
struct segment {
  int i, l, r;
  segment() {}
  segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
  dp[0] = 0;
  deque<segment> dq; dq.push_back(segment(0, 1, n));
for (int i = 1; i <= n; ++i) {</pre>
    dp[i] = f(dq.front().i, i);
    while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
    dq.front().l = i + 1
    segment seg = segment(i, i + 1, n);
    while (dq.size() &&
      f(i, dq.back().1) < f(dq.back().i, dq.back().1))
         dq.pop_back();
    if (dq.size()) {
      int d = 1 << 20, c = dq.back().1;</pre>
      while (d >>= 1) if (c + d <= dq.back().r)</pre>
         if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
```

```
dq.back().r = c; seg.1 = c + 1;
}
if (seg.1 <= n) dq.push_back(seg);
}
}</pre>
```

8.5 ConvexHull Optimization

```
inline lld DivCeil(lld n, lld d) { // ceil(n/d)
  return n / d + (((n < 0) != (d > 0)) && (n % d));
struct Line {
  static bool flag;
  lld a, b, l, r; // y=ax+b in [l, r)
  11d operator()(11d x) const { return a * x + b; }
  bool operator<(const Line& i) const {</pre>
    return flag ? tie(a, b) < tie(i.a, i.b) : 1 < i.l;</pre>
  11d operator&(const Line& i) const {
    return DivCeil(b - i.b, i.a - a);
bool Line::flag = true;
class ConvexHullMax {
  set<Line> L;
 public:
  ConvexHullMax() { Line::flag = true; }
  void InsertLine(lld a, lld b) { // add y = ax + b
    Line now = \{a, b, -INF, INF\};
    if (L.empty())
      L.insert(now);
      return;
    Line::flag = true;
    auto it = L.lower_bound(now);
    auto prv = it == L.begin() ? it : prev(it);
    if (it != L.end() && ((it != L.begin() &&
      (*it)(it->1) >= now(it->1) &&
      (*prv)(prv->r - 1) >= now(prv->r - 1)) ||
      (it == L.begin() && it->a == now.a))) return;
    if (it != L.begin()) {
      while (prv != L.begin() &&
        (*prv)(prv->1) \le now(prv->1))
                 --L.erase(prv);
      if (prv == L.begin() && now.a == prv->a)
        L.erase(prv);
    if (it != L.end())
      while (it != --L.end() &&
        (*it)(it->r) <= now(it->r))
          it = L.erase(it);
    if (it != L.begin()) {
      prv = prev(it);
      const_cast<Line*>(&*prv)->r=now.l=((*prv)&now);
    if (it != L.end())
      const_cast<Line*>(&*it)->l=now.r=((*it)&now);
    L.insert(it, now);
  11d Query(11d a) const { // query max at x=a
    if (L.empty()) return -INF;
    Line::flag = false;
    auto it = --L.upper_bound({0, 0, a, 0});
    return (*it)(a);
};
```

8.6 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
      s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
```

8.7 Cactus Matching

}

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
  dfn[u]=low[u]=++dfs_idx;
  for(int i=0;i<(int)init_g[u].size();i++){</pre>
    int v=init_g[u][i];
    if(v==par[u]) continue;
    if(!dfn[v]){
      par[v]=u;
      tarjan(v);
      low[u]=min(low[u],low[v]);
      if(dfn[u]<low[v]){</pre>
        g[u].push_back(v);
        g[v].push_back(u);
    }else{
      low[u]=min(low[u],dfn[v]);
      if(dfn[v]<dfn[u]){</pre>
        int temp_v=u;
        bcc_id++;
        while(temp_v!=v){
          g[bcc_id+n].push_back(temp_v);
          g[temp_v].push_back(bcc_id+n);
           temp_v=par[temp_v];
        g[bcc_id+n].push_back(v);
        g[v].push_back(bcc_id+n);
        reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
    }
  }
int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
void dfs(int u,int fa){
  if(u<=n){
    for(int i=0;i<(int)g[u].size();i++){</pre>
      int v=g[u][i];
      if(v==fa) continue;
      dfs(v,u);
      memset(tp,0x8f,sizeof tp);
      if(v<=n){
        tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
        tp[1]=max(
          dp[u][0]+dp[v][0]+1
          dp[u][1]+max(dp[v][0],dp[v][1])
      }else{
        tp[0]=dp[u][0]+dp[v][0];
        tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
      dp[u][0]=tp[0],dp[u][1]=tp[1];
    }
  }else{
    for(int i=0;i<(int)g[u].size();i++){</pre>
      int v=g[u][i];
      if(v==fa) continue;
      dfs(v,u);
    min_dp[0][0]=0;
    min_dp[1][1]=1;
    min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
    for(int i=0;i<(int)g[u].size();i++){</pre>
      int v=g[u][i];
      if(v==fa) continue;
      memset(tmp,0x8f,sizeof tmp);
      tmp[0][0]=max(
        \min_{dp[0][0]+\max(dp[v][0],dp[v][1])}
        min_dp[0][1]+dp[v][0]
      tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
      tmp[1][0]=max
        \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
        min_dp[1][1]+dp[v][0]
      tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
      memcpy(min_dp,tmp,sizeof tmp);
```

ans[d] = row[i];

for(int j = R[i]; j != i; j = R[j])

```
dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
                                                                          remove(col[j]);
    dp[u][0]=min_dp[0][0];
                                                                        dance(d+1);
                                                                        for(int j = L[i]; j != i; j = L[j])
                                                                          resume(col[j]);
int main(){
 int m,a,b;
                                                                      resume(c);
  scanf("%d%d",&n,&m);
                                                                   }
  for(int i=0;i<m;i++){</pre>
                                                                 } sol;
    scanf("%d%d",&a,&b);
    init_g[a].push_back(b);
                                                                 8.9 Tree Knapsack
    init_g[b].push_back(a);
                                                                 int dp[N][K];PII obj[N];
 par[1]=-1;
                                                                 vector<int> G[N];
 tarjan(1);
                                                                 void dfs(int u, int mx){
 dfs(1,-1);
                                                                   for(int s: G[u])
 printf("%d\n", max(dp[1][0], dp[1][1]));
                                                                      if(mx < obj[s].first) continue;</pre>
 return 0:
                                                                      for(int i=0;i<=mx-obj[s].FF;i++)</pre>
                                                                        dp[s][i] = dp[u][i];
                                                                      dfs(s, mx - obj[s].first);
                                                                      for(int i=obj[s].FF;i<=mx;i++)</pre>
8.8
      DLX
                                                                        dp[u][i] = max(dp[u][i],
                                                                          dp[s][i - obj[s].FF] + obj[s].SS);
struct DLX {
                                                                   }
  const static int maxn=210;
  const static int maxm=210;
                                                                 int main(){
  const static int maxnode=210*210;
                                                                   int n, k; cin >> n >> k;
 int n, m, size, row[maxnode], col[maxnode];
int U[maxnode], D[maxnode], L[maxnode], R[maxnode];
                                                                   for(int i=1;i<=n;i++){</pre>
                                                                     int p; cin >> p;
  int H[maxn], S[maxm], ansd, ans[maxn];
                                                                     G[p].push_back(i);
 void init(int _n, int _m) {
    n = _n, m = _m;
                                                                     cin >> obj[i].FF >> obj[i].SS;
    for(int i = 0; i <= m; ++i) {</pre>
                                                                   dfs(0, k); int ans = 0;
      S[i] = 0;
                                                                   for(int i=0;i<=k;i++) ans = max(ans, dp[0][i]);</pre>
      U[i] = D[i] = i;
                                                                   cout << ans << '\n';
      L[i] = i-1, R[i] = i+1;
                                                                   return 0;
    R[L[0] = size = m] = 0;
    for(int i = 1; i <= n; ++i) H[i] = -1;
                                                                 8.10
                                                                          N Queens Problem
  void Link(int r, int c) {
    ++S[col[++size] = c];
                                                                 vector< int > solve( int n ) {
    row[size] = r; D[size] = D[c];
                                                                   // no solution when n=2, 3
    U[D[c]] = size; U[size] = c; D[c] = size;
if(H[r] < 0) H[r] = L[size] = R[size] = size;</pre>
                                                                   vector< int > ret;
                                                                   if ( n % 6 == 2 ) {
  for ( int i = 2 ; i <= n ; i += 2 )</pre>
    else {
      R[size] = R[H[r]];
                                                                        ret.push_back( i );
      L[R[H[r]]] = size;
                                                                      ret.push_back( 3 ); ret.push_back( 1 );
for ( int i = 7 ; i <= n ; i += 2 )
      L[size] = H[r];
      R[H[r]] = size;
                                                                        ret.push_back( i );
                                                                      ret.push_back( 5 );
                                                                   } else if ( n % 6 == 3 ) {
  void remove(int c) {
                                                                      for ( int i = 4 ; i <= n ; i += 2 )
    L[R[c]] = L[c]; R[L[c]] = R[c];
                                                                        ret.push_back( i );
    for(int i = D[c]; i != c; i = D[i])
                                                                      ret.push_back( 2 );
      for(int j = R[i]; j != i; j = R[j]) {
                                                                      for ( int i = 5 ; i <= n ; i += 2 )
        U[D[j]] = U[j];
                                                                        ret.push_back( i );
        D[U[j]] = D[j];
                                                                      ret.push_back( 1 ); ret.push_back( 3 );
         --S[col[j]];
                                                                   } else {
                                                                      for ( int i = 2 ; i <= n ; i += 2 )
                                                                      ret.push_back( i );
for ( int i = 1 ; i <= n ; i += 2 )
  void resume(int c) {
    L[R[c]] = c; R[L[c]] = c;
for(int i = U[c]; i != c; i = U[i])
                                                                        ret.push_back( i );
      for(int j = L[i]; j != i; j = L[j]) {
                                                                    return ret;
        U[D[j]] = j;
D[U[j]] = j;
        ++S[col[j]];
    }
  void dance(int d) {
    if(d>=ansd) return;
    if(R[0] == 0) {
      ansd = d;
      return;
    int c = R[0];
    for(int i = R[0]; i; i = R[i])
      if(S[i] < S[c]) c = i;
    remove(c);
    for(int i = D[c]; i != c; i = D[i]) {
```