# **Contents**

Contents				
1 B	asic	1	5.25 Simplex Construction	15
1.1	vimrc	i	5.26 Adaptive Simpson	15
1.2	Debug Macro	1	•	
1.3	Increase Stack	1	6 Geometry	15
1.4	Pragma Optimization	1	6.1 Basic Geometry	15
1.5	IO Optimization	1	6.2 2D Convex Hull	15
1.6	SVG Writer	1	6.3 2D Farthest Pair	15
2 D	ata Structure	2	6.4 MinMax Enclosing Rect	15
2.1	Dark Magic	2	6.5 Minkowski Sum	16
2.2	Link-Cut Tree	2	6.6 Segment Intersection	16
2.3	LiChao Segment Tree	2	6.7 Half Plane Intersec-	
2.4	Treap	3	tion	16
2.5	Linear Basis	3	6.8 SegmentDist	
2.6	Binary Search On		(Sausage)	16
	Segtree	3	6.9 Rotating Sweep Line	16
3 G	raph	3	6.10 Polygon Cut	16
3.1		3	6.11 Point In Simple	
3.1	2-SAT (SCC)	3	Polygon	17
3.3	Round Square Tree	4	6.12 Point In Hull (Fast)	17
3.4	Edge TCC	4	6.13 Point In Polygon	
3.5	DMST	4	(Fast)	17
3.6	Dominator Tree	4	6.14 Tangent of Points	
3.7	Edge Coloring	5	To Hull	17
3.8	Centroid Decom-		6.15 Circle Class & Intersection	17
0.0	position	5		17
3.9	Lowbit Decompo-		6.16 Circle Common Tangent	17
	sition	5	6.17 Line-Circle Inter-	
	Virtual Tree	6	section	18
3.11	Tree Hashing	6	6.18 Poly-Circle Inter-	18
3.12	Mo's Algorithm on	_	section	10
3.13	Tree	6 6	Circle	18
		6	6.20 Circle Union	18
	MaximalClique		6.21 Polygon Union	18
3.15	MaximumClique	6	6.22 3D Point	19
3.16	Minimum Mean Cycle	7	6.23 3D Convex Hull	19
, FI			6.24 3D Projection	19
	ow & Matching	7	6.25 3D Skew Line Near-	
4.1	HopcroftKarp	7	est Point	19 19
4.2	Dijkstra Cost Flow	7	6.26 Delaunay	
4.3	<u>Dinic</u>	8	6.27 Build Voronoi 6.28 kd Tree (Nearest	20
4.4 4.5	Flow Models	8	Point)	20
4.5	General Graph Matching	8	6.29 kd Closest Pair (3D	
4.6	Global Min-Cut	9	ver.)	20
4.7	GomoryHu Tree	9	6.30 Simulated Annealing	20
4.8	Kuhn Munkres	9	6.31 Triangle Centers	20
4.9	Minimum Cost Cir-		_	
	culation	9	7 Stringology	21
4.10	Minimum Cost Max Flow	9	7.1 Hash	21
4.11	Weighted Matching .	10	7.2 Suffix Array	21
	•		7.3 Ex SAM	21 21
<b>5 M</b> 5.1		11	7.5 Manacher	21
5.1 5.2	Common Bounds Stirling Number	11 11	7.6 Lyndon Factorization	22
5.2 5.3	•		7.7 Main Lorentz	22
	ax+by=gcd	11	7.8 BWT	22
5.4 5.5	Chinese Remainder .	11 11	7.9 Palindromic Tree	22
	DiscreteLog		8 Misc	22
5.6 5.7	Quadratic Residue	11	8.1 Theorems	22
5.7 5.8	Extended Euler Extended FloorSum .	11 11	8.2 Weight Matroid In-	
5.9	Extended Euclidean .	12	tersection	23
5.10		12	8.3 Stable Marriage	23
5.11	ModMin	12	8.4 Bitset LCS	23
5.12	FWT	12	8.5 Prefix Substring LCS.	23
5.13 5.14	Packed FFT CRT for arbitrary mod	12 12	8.6 Convex 1D/1D DP	23
			8.7 ConvexHull Opti-	~-
5.15 5.16	NTT / FFT Formal Power Series	12 13	mization	23
5.16		13	8.8 Min Plus Convolution 8.9 De-Bruijn	24 24
5.17		i.J	-	24 24
5.10	Sieve)	13	8.10 Josephus Problem	
5.19		14	8.11 N Queens Problem	24
5.20	Pollard Rho	14	8.12 Tree Knapsack	24
	Berlekamp Massey	14	8.13 Manhattan MST	24
	Gauss Elimination	14	8.14 Binary Search On Fraction	24
5.23	Charateristic Poly-	٦,	8.15 Barrett Reduction	24
52/	nomial	14 14	8.16 Montgomery	24

# 1 Basic

# 1.2 Debug Macro [851d50]

```
#define all(x) begin(x), end(x)
#ifdef CKISEKI
#include <experimental/iterator>
#define safe cerr<<__PRETTY_FUNCTION__<<" line "<</pre>
      __LINE__<<" safe\n"
#define debug(a...) debug_(#a, a)
#define orange(a...) orange_(#a, a)
void debug_(const char *s, auto ...a) {
  cerr << "\e[1;32m(" << s << ") = (";</pre>
 int f = 0;
 (..., (cerr << (f++ ? ", " : "") << a));
 cerr << ")\e[0m\n";</pre>
void orange_(const char *s, auto L, auto R) {
  cerr << "\e[1;33m[ " << s << " ] = [ ";</pre>
 using namespace experimental;
 copy(L, R, make_ostream_joiner(cerr, ", "));
cerr << " ]\e[0m\n";</pre>
#else
#define safe ((void)0)
#define debug(...) safe
#define orange(...) safe
#endif
```

#### 1.3 Increase Stack [b6856c]

```
const int size = 256 << 20;
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));</pre>
```

#### 1.4 Pragma Optimization [6006f6]

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,arch=skylake")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
```

#### 

```
static inline int gc() {
  constexpr int B = 1<<20; static char buf[B], *p, *q;
  if (p == q) q = (p = buf) + fread(buf, 1, B, stdin);
  return q == buf ? EOF : *p++;
}</pre>
```

#### 1.6 SVG Writer [57436c]

```
class SVG {
  void p(string_view s) { o << s; }
  void p(string_view s, auto v, auto... vs) {
    auto i = s.find('$');
    o << s.substr(0, i) << v, p(s.substr(i + 1), vs...);
  }
  ofstream o; string c = "red";
public:
  SVG(auto f, auto x1, auto y1, auto x2, auto y2) : o(f) {
    p("<svg xmlns='http://www.w3.org/2000/svg' "
        "viewBox='$ $ $'>\n"
        "<style>*{stroke-width:0.5%;}</style>\n",
        x1, -y2, x2 - x1, y2 - y1); }
```

```
~SVG() { p("</svg>\n"); }
SVG &color(string nc) { return c = nc, *this; }
void line(auto x1, auto y1, auto x2, auto y2) {
p("<line x1='$' y1='$' x2='$' y2='$' stroke='$'/>\n",
  x1, -y1, x2, -y2, c); }
void circle(auto x, auto y, auto r) {
p("<circle cx='$' cy='$' r='$' stroke='$' "</pre>
   "fill='none'/>\n", x, -y, r, c); }
void text(auto x, auto y, string s, int w = 12) {
 p("<text x='$' y='$' font-size='$px'>$</text>\n",
  x, -y, w, s); }
```

#### 2 **Data Structure**

# 2.1 Dark Magic [095f25]

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using namespace __gnu_pbds;
// heap tags: paring/binary/binomial/rc_binomial/thin
template<typename T>
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>, \
                  pairing_heap_tag>;
// pbds_heap::point_iterator
  ' x = pq.push(10); pq.modify(x, 87); a.join(b);
// tree tags: rb_tree_tag/ov_tree_tag/splay_tree_tag
template<typename T>
using ordered_set = tree<T, null_type, less<T>,
  rb_tree_tag, tree_order_statistics_node_update>;
// find_by_order, order_of_key
// hash tables: cc_hash_table/gp_hash_table
```

#### 2.2 Link-Cut Tree [5e4f69]

```
template <typename Val, typename SVal> class LCT {
struct node {
 int pa, ch[2];
 bool rev;
 Val v, prod, rprod;
  SVal sv, sub, vir;
 node() : pa{0}, ch{0, 0}, rev{false}, v{}, prod{},
    rprod{}, sv{}, sub{}, vir{} {};
#define cur o[u]
#define lc cur.ch[0]
#define rc cur.ch[1]
vector<node> o;
bool is_root(int u) const {
 return o[cur.pa].ch[0] != u && o[cur.pa].ch[1] != u;
bool is_rch(int u) const {
 return o[cur.pa].ch[1] == u && !is_root(u); }
 void down(int u) {
  if (not cur.rev) return;
 if (lc) set_rev(lc);
 if (rc) set_rev(rc);
 cur.rev = false;
void up(int u) {
 cur.prod = o[lc].prod * cur.v * o[rc].prod;
  cur.rprod = o[rc].rprod * cur.v * o[lc].rprod;
 cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
void set_rev(int u) {
 swap(lc, rc), swap(cur.prod, cur.rprod);
cur.rev ^= 1;
void rotate(int u) {
 int f = cur.pa, g = o[f].pa, l = is_rch(u);
if (cur.ch[l ^ 1]) o[cur.ch[l ^ 1]].pa = f;
 if (not is_root(f)) o[g].ch[is_rch(f)] = u;
 o[f].ch[l] = cur.ch[l ^ 1];
 cur.ch[l ^ 1] = f;
 cur.pa = g, o[f].pa = u;
 up(f);
}
void splay(int u) {
 vector<int> stk = {u};
 while (not is_root(stk.back()))
   stk.push_back(o[stk.back()].pa);
 while (not stk.empty()) {
  down(stk.back());
   stk.pop_back();
 }
```

```
for (int f = cur.pa; not is_root(u); f = cur.pa) {
   if (!is_root(f))
    rotate(is_rch(u) == is_rch(f) ? f : u);
   rotate(u);
  up(u);
 void access(int x) {
  for (int u = x, last = 0; u; u = cur.pa) {
   splay(u);
   cur.vir = cur.vir + o[rc].sub - o[last].sub;
   rc = last; up(last = u);
  splay(x);
 int find_root(int u) {
  int la = 0;
  for (access(u); u; u = lc) down(la = u);
  return la;
 void split(int x, int y) { chroot(x); access(y); }
 void chroot(int u) { access(u); set_rev(u); }
public:
 LCT(int n = 0) : o(n + 1) {}
 int add(const Val &v = {}) {
  o.push_back(v);
  return int(o.size()) - 2;
 void set_val(int u, const Val &v) {
  splay(++u); cur.v = v; up(u);
 void set_sval(int u, const SVal &v) {
  splay(++u); cur.sv = v; up(u);
 Val query(int x, int y) {
  split(++x, ++y); return o[y].prod;
 SVal subtree(int p, int u) {
  chroot(++p); access(++u);
  return cur.vir + cur.sv;
 bool connected(int u, int v) {
  return find_root(++u) == find_root(++v); }
 void link(int x, int y) {
  chroot(++x); access(++y);
  o[y].vir = o[y].vir + o[x].sub;
  up(o[x].pa = y);
 void cut(int x, int y) {
  split(++x, ++y);
  o[y].ch[0] = o[x].pa = 0; up(y);
#undef cur
#undef lc
#undef rc
};
2.3 LiChao Segment Tree [b9c827]
```

```
struct L {
 int m, k, id;
 L() : id(-1) {}
 L(int a, int b, int c) : m(a), k(b), id(c) {}
 int at(int x) { return m * x + k; }
};
class LiChao {
private:
 int n; vector<L> nodes;
 static int lc(int x) { return 2 * x + 1; }
static int rc(int x) { return 2 * x + 2; }
 void insert(int l, int r, int id, L ln) {
  int m = (l + r) >> 1;
  if (nodes[id].id == -1)
   return nodes[id] = ln, void();
  bool atLeft = nodes[id].at(l) < ln.at(l);</pre>
  if (nodes[id].at(m) < ln.at(m))</pre>
   atLeft ^= 1, swap(nodes[id], ln);
  if (r - l == 1) return;
  if (atLeft) insert(l, m, lc(id), ln);
  else insert(m, r, rc(id), ln);
 int query(int l, int r, int id, int x) {
  int m = (l + r) >> 1, ret = 0;
```

```
National Taiwan University - ckiseki
  if (nodes[id].id != -1) ret = nodes[id].at(x);
  if (r - l == 1) return ret;
  if (x < m) return max(ret, query(l, m, lc(id), x));</pre>
  return max(ret, query(m, r, rc(id), x));
public:
LiChao(int n_{-}): n(n_{-}), nodes(n * 4) {}
void insert(L ln) { insert(0, n, 0, ln); }
int query(int x) { return query(0, n, 0, x); }
};
     Treap [ae576c]
 _gnu_cxx::sfmt19937 rnd(7122); // <ext/random>
namespace Treap {
struct node {
int size, pri; node *lc, *rc, *pa;
node() : size(1), pri(rnd()), lc(0), rc(0), pa(0) {}
void pull() {
 size = 1; pa = 0;
  if (lc) { size += lc->size; lc->pa = this; }
```

```
if (rc) { size += rc->size; rc->pa = this; }
int SZ(node *x) { return x ? x->size : 0; }
node *merge(node *L, node *R) {
if (not L or not R) return L ? L : R;
if (L->pri > R->pri)
 return L->rc = merge(L->rc, R), L->pull(), L;
else
 return R->lc = merge(L, R->lc), R->pull(), R;
void splitBySize(node *o, int k, node *&L, node *&R) {
if (not o) L = R = 0;
else if (int s = SZ(o\rightarrow lc) + 1; s \le k)
 L=o, splitBySize(o->rc, k-s, L->rc, R), L->pull();
else
 R=o, splitBySize(o->lc, k, L, R->lc), R->pull();
 // SZ(L) == k
int getRank(node *o) { // 1-base
```

if (o->pa->rc == o) r += SZ(o->pa->lc) + 1;

# 2.5 Linear Basis [138d5d]

**int** r = SZ(o->lc) + 1;

} // namespace Treap

return r;

**for** (; o->pa; o = o->pa)

```
template <int BITS, typename S = int> struct Basis {
 static constexpr S MIN = numeric_limits<S>::min();
 array<pair<llu, S>, BITS> b;
 Basis() { b.fill({0, MIN}); }
 void add(llu x, S p) {
  for (int i = BITS-1; i>=0; i--) if (x >> i & 1) {
   if (b[i].first == 0) return b[i]={x, p}, void();
   if (b[i].second < p)</pre>
    swap(b[i].first, x), swap(b[i].second, p);
   x ^= b[i].first;
  }
 optional<llu> query_kth(llu v, llu k) {
  vector<pair<llu, int>> o;
for (int i = 0; i < BITS; i++)</pre>
   if (b[i].first) o.emplace_back(b[i].first, i);
  if (k >= (1ULL << o.size())) return {};</pre>
  for (int i = int(o.size()) - 1; i >= 0; i--)
   if ((k >> i & 1) ^ (v >> o[i].second & 1))
    v ^= o[i].first;
  return v;
 Basis filter(S l) {
  Basis res = *this;
  for (int i = 0; i < BITS; i++)</pre>
   if (res.b[i].second < l) res.b[i] = {0, MIN};</pre>
  return res;
};
```

# 2.6 Binary Search On Segtree [6c61c0]

```
// find_first = l -> minimal x s.t. check( [l, x) )
// find_last = r -> maximal x s.t. check( [x, r) )
int find_first(int l, auto &&check) {
  if (l >= n) return n + 1;
```

```
l += sz; push(l); Monoid sum; // identity
 while ((l & 1) == 0) l >>= 1;
 if (auto s = sum + nd[l]; check(s)) {
  while (l < sz) {</pre>
   prop(l); l = (l << 1);
    if (auto nxt = sum + nd[l]; not check(nxt))
     sum = nxt, l++;
  return l + 1 - sz;
 } else sum = s, l++;
} while (lowbit(l) != l);
return n + 1;
int find_last(int r, auto &&check) {
if (r <= 0) return -1;
r += sz; push(r - 1); Monoid sum; // identity
 while (r > 1 and (r & 1)) r >>= 1;
 if (auto s = nd[r] + sum; check(s)) {
  while (r < sz) {
   prop(r); r = (r << 1) | 1;
    if (auto nxt = nd[r] + sum; not check(nxt))
     sum = nxt, r--;
  return r - sz;
 } else sum = s;
} while (lowbit(r) != r);
return -1;
```

# 3 Graph

# 3.1 2-SAT (SCC) [09167a]

```
class TwoSat { // test @ CSES Giant Pizza
private:
 int n; vector<vector<int>>> G, rG, sccs;
 vector<int> ord, idx, vis, res;
 void dfs(int u) {
  vis[u] = true;
  for (int v : G[u]) if (!vis[v]) dfs(v);
  ord.push_back(u);
 void rdfs(int u) {
  vis[u] = false; idx[u] = sccs.size() - 1;
  sccs.back().push_back(u);
  for (int v : rG[u]) if (vis[v]) rdfs(v);
public:
 TwoSat(int n_{-}): n(n_{-}), G(n), rG(n), idx(n), vis(n),
    res(n) {}
 void add_edge(int u, int v) {
  G[u].push_back(v); rG[v].push_back(u);
 void orr(int x, int y) {
  if ((x ^ y) == 1) return;
  add_edge(x ^ 1, y); add_edge(y ^ 1, x);
 bool solve() {
  for (int i = 0; i < n; ++i) if (not vis[i]) dfs(i);</pre>
  for (int u : ord | views::reverse)
   if (vis[u]) sccs.emplace_back(), rdfs(u);
  for (int i = 0; i < n; i += 2)</pre>
   if (idx[i] == idx[i + 1]) return false;
  vector<bool> c(sccs.size());
  for (size_t i = 0; i < sccs.size(); ++i)</pre>
   for (int z : sccs[i])
    res[z] = c[i], c[idx[z ^ 1]] = !c[i];
  return true;
 bool get(int x) { return res[x]; }
 int get_id(int x) { return idx[x]; }
 int count() { return sccs.size(); }
};
```

#### **3.2** BCC [6ac6db]

```
class BCC {
  int n, ecnt, bcnt;
  vector<vector<pair<int, int>>> g;
  vector<int> dfn, low, bcc, stk;
  vector<bool> ap, bridge;
  void dfs(int u, int f) {
```

```
dfn[u] = low[u] = dfn[f] + 1;
  int ch = 0;
  for (auto [v, t] : g[u]) if (bcc[t] == -1) {
   bcc[t] = 0; stk.push_back(t);
   if (dfn[v]) {
    low[u] = min(low[u], dfn[v]);
    continue;
   ++ch, dfs(v, u);
  low[u] = min(low[u], low[v]);
if (low[v] > dfn[u]) bridge[t] = true;
   if (low[v] < dfn[u]) continue;</pre>
   ap[u] = true;
   while (not stk.empty()) {
    int o = stk.back(); stk.pop_back();
    bcc[o] = bcnt;
    if (o == t) break;
  bcnt += 1;
 ap[u] = ap[u] and (ch != 1 or u != f);
public:
BCC(int n_{-}) : n(n_{-}), ecnt(0), bcnt(0), g(n), dfn(n),
    low(n), stk(), ap(n) {}
void add_edge(int u, int v) {
 g[u].emplace_back(v, ecnt);
  g[v].emplace_back(u, ecnt++);
void solve() {
 bridge.assign(ecnt, false); bcc.assign(ecnt, -1);
 for (int i = 0; i < n; ++i) if (!dfn[i]) dfs(i, i);</pre>
int bcc_id(int x) const { return bcc[x]; }
bool is_ap(int x) const { return ap[x]; }
bool is_bridge(int x) const { return bridge[x]; }
```

# 3.3 Round Square Tree [528440]

```
struct RST {
 int n; vector<vector<int>> T;
 RST(auto &G) : n(G.size()), T(n) {
  vector<int> stk, vis(n), low(n);
  auto dfs = [&](auto self, int u, int d) -> void {
   low[u] = vis[u] = d; stk.push_back(u);
   for (int v : G[u]) if (!vis[v]) {
    self(self, v, d + 1);
if (low[v] == vis[u]) {
     int cnt = T.size(); T.emplace_back();
     for (int x = -1; x != v; stk.pop_back())
      T[cnt].push_back(x = stk.back());
     T[u].push_back(cnt); // T is rooted
    } else low[u] = min(low[u], low[v]);
   } else low[u] = min(low[u], vis[v]);
  };
  for (int u = 0; u < N; u++)</pre>
   if (!vis[u]) dfs(dfs, u, 1);
} // T may be forest; after dfs, stk are the roots
}; // test @ 2020 Shanghai K
```

#### 3.4 Edge TCC [5a2668]

```
vector<vector<int>> ETCC(auto &adj) {
const int n = static_cast<int>(adj.size());
vector<int> up(n), low(n), in, out, nx, id;
 in = out = nx = id = vector < int > (n, -1);
int dfc = 0, cnt = 0; Dsu dsu(n);
auto merge = [&](int u, int v) {
  dsu.join(u, v); up[u] += up[v]; };
auto dfs = [&](auto self, int u, int p) -> void {
  in[u] = low[u] = dfc++;
  for (int v : adj[u]) if (v != u) {
   if (v == p) { p = -1; continue; }
   if (in[v] == -1) {
    self(self, v, u);
if (nx[v] == -1 && up[v] <= 1) {</pre>
     up[u] += up[v]; low[u] = min(low[u], low[v]);
     continue;
    if (up[v] == 0) v = nx[v];
    if (low[u] > low[v])
     low[u] = low[v], swap(nx[u], v);
    for (; v != -1; v = nx[v]) merge(u, v);
```

```
} else if (in[v] < in[u]) {</pre>
    low[u] = min(low[u], in[v]); up[u]++;
   } else {
    for (int &x = nx[u]; x != -1 &&
      in[x] <= in[v] && in[v] < out[x]; x = nx[x])
     merge(u, x);
    up[u]--;
 out[u] = dfc;
 for (int i = 0; i < n; i++)</pre>
  if (in[i] == -1) dfs(dfs, i, -1);
 for (int i = 0; i < n; i++)</pre>
  if (dsu.anc(i) == i) id[i] = cnt++;
 vector<vector<int>> comps(cnt);
 for (int i = 0; i < n; i++)</pre>
 comps[id[dsu.anc(i)]].push_back(i);
 return comps;
} // test @ yosupo judge
```

#### **3.5** DMST [75c30d]

```
using D = int64_t;
struct E { int s, t; D w; }; // 0-base
vector<int> dmst(const vector<E> &e, int n, int root) {
 using PQ = pair<min_heap<pair<D, int>>, D>;
 auto push = [](PQ &pq, pair<D, int> v) {
 pq.first.emplace(v.first - pq.second, v.second); };
 auto top = [](const PQ &pq) -> pair<D, int> {
  auto r = pq.first.top();
  return {r.first + pq.second, r.second}; };
 auto join = [&push, &top](PQ &a, PQ &b) {
  if (a.first.size() < b.first.size()) swap(a, b);</pre>
  for (; !b.first.empty(); b.first.pop())
   push(a, top(b)); };
 vector<PQ> h(n * 2);
 for (size_t i = 0; i < e.size(); ++i)</pre>
 push(h[e[i].t], {e[i].w, i});
 vector<int> a(n*2), v(n*2, -1), pa(n*2, -1), r(n*2);
 iota(a.begin(), a.end(), 0);
 auto o = [&](int x) { int y;
  for (y = x; a[y] != y; y = a[y]);
  for (int ox = x; x != y; ox = x)
  x = a[x], a[ox] = y;
  return y; };
 int pc = (v[root] = n + 1) - 1;
 for (int i = 0; i < n; ++i) if (v[i] == -1)
  for (int p=i; v[p]<0||v[p]==i; p=o(e[r[p]].s)) {</pre>
   if (int q = p; v[q] == i && (p = pc++, 1)) do {
    h[q].second = -h[q].first.top().first;
    join(h[pa[q] = a[q] = p], h[q]);
   } while ((q = o(e[r[q]].s)) != p);
   \textbf{for}(v[p] = i; !h[p].first.empty()\&\&o(e[top(h[p]).second
    ].s)==p;h[p].first.pop());
   r[p] = top(h[p]).second;
 vector<int> ans;
 for (int i=pc-1;i>=0;i--) if (i!=root&&v[i]!=n) {
 for (int f = e[r[i]].t; f!=-1&&v[f]!=n; f = pa[f])
   v[f] = n;
 ans.push_back(r[i]);
 return ans; // default minimize, returns edgeid array
```

#### 3.6 Dominator Tree [ea5b7c]

```
struct Dominator {
 vector<vector<int>> g, r, rdom; int tk;
vector<int> dfn, rev, fa, sdom, dom, val, rp;
 Dominator(int n) : g(n), r(n), rdom(n), tk(0) {
 dfn = rev = fa = sdom = dom =
   val = rp = vector<int>(n, -1); }
 void add_edge(int x, int y) { g[x].push_back(y); }
 void dfs(int x) {
  rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk; tk++;
  for (int u : g[x]) {
   if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
   r[dfn[u]].push_back(dfn[x]);
  }
 void merge(int x, int y) { fa[x] = y; }
```

```
int find(int x, int c = 0) {
  if (fa[x] == x) return c ? -1 : x;
 if (int p = find(fa[x], 1); p != -1) {
   if (sdom[val[x]] > sdom[val[fa[x]]])
   val[x] = val[fa[x]];
  fa[x] = p;
   return c ? p : val[x];
 } else return c ? fa[x] : val[x];
vector<int> build(int s, int n) {
// return the father of each node in dominator tree
 dfs(s); // p[i] = -2 if i is unreachable from s
 for (int i = tk - 1; i >= 0; --i) {
  for (int u : r[i])
   sdom[i] = min(sdom[i], sdom[find(u)]);
   if (i) rdom[sdom[i]].push_back(i);
   for (int u : rdom[i]) {
   int p = find(u);
   dom[u] = (sdom[p] == i ? i : p);
  if (i) merge(i, rp[i]);
 }
 vector<int> p(n, -2); p[s] = -1;
 for (int i = 1; i < tk; ++i)</pre>
   if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
 for (int i = 1; i < tk; ++i)
  p[rev[i]] = rev[dom[i]];
 return p;
} // test @ yosupo judge
     Edge Coloring [029763]
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
```

```
void clear(int N) {
for (int i = 0; i <= N; i++)
for (int j = 0; j <= N; j++)</pre>
    C[i][j] = G[i][j] = 0;
void solve(vector<pair<int, int>> &E, int N) {
int X[kN] = {}, a;
auto update = [&](int u) {
 for (X[u] = 1; C[u][X[u]]; X[u]++);
auto color = [&](int u, int v, int c) {
 int p = G[u][v];
 G[u][v] = G[v][u] = c;
 C[u][c] = v, C[v][c] = u;
 C[u][p] = C[v][p] = 0;
 if (p) X[u] = X[v] = p;
 else update(u), update(v);
 return p;
};
auto flip = [&](int u, int c1, int c2) {
 int p = C[u][c1];
 swap(C[u][c1], C[u][c2]);
 if (p) G[u][p] = G[p][u] = c2;
 if (!C[u][c1]) X[u] = c1;
 if (!C[u][c2]) X[u] = c2;
 return p;
};
for (int i = 1; i <= N; i++) X[i] = 1;</pre>
for (int t = 0; t < E.size(); t++) {</pre>
 auto [u, v] = E[t];
  int v0 = v, c = X[u], c0 = c, d;
 vector<pair<int, int>> L; int vst[kN] = {};
 while (!G[u][v0]) {
   L.emplace_back(v, d = X[v]);
  if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
     c = color(u, L[a].first, c);
  else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
    color(u, L[a].first, L[a].second);
   else if (vst[d]) break;
  else vst[d] = 1, v = C[u][d];
 if (!G[u][v0]) {
   for (; v; v = flip(v, c, d), swap(c, d));
  if (C[u][c0]) { a = int(L.size()) - 1;
   while (--a >= 0 && L[a].second != c);
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
  } else t--;
```

```
3.8 Centroid Decomposition [63b2fb]
```

```
struct Centroid {
 using G = vector<vector<pair<int, int>>>;
 vector<vector<int64_t>> Dist;
 vector<int> Pa, Dep;
 vector<int64_t> Sub, Sub2;
 vector<int> Cnt, Cnt2;
 vector<int> vis, sz, mx, tmp;
 void DfsSz(const G &g, int x) {
  vis[x] = true, sz[x] = 1, mx[x] = 0;
for (auto [u, w] : g[x]) if (not vis[u]) {
   DfsSz(g, u); sz[x] += sz[u];
   mx[x] = max(mx[x], sz[u]);
  tmp.push_back(x);
 void DfsDist(const G &g, int x, int64_t D = 0) {
  Dist[x].push_back(D); vis[x] = true;
  for (auto [u, w] : g[x])
   if (not vis[u]) DfsDist(g, u, D + w);
 void DfsCen(const G &g, int x, int D = 0, int p = -1)
  tmp.clear(); DfsSz(g, x);
  int M = tmp.size(), C = -1;
  for (int u : tmp) {
   if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
   vis[u] = false;
  DfsDist(g, C);
  for (int u : tmp) vis[u] = false;
  Pa[C] = p, vis[C] = true, Dep[C] = D;
  for (auto [u, w] : g[C])
   if (not vis[u]) DfsCen(g, u, D + 1, C);
 Centroid(int N, G g)
   : Sub(N), Sub2(N), Cnt(N), Cnt2(N), Dist(N), Pa(N),
    Dep(N), vis(N), sz(N), mx(N) { DfsCen(g, 0); }
 void Mark(int v) {
  int x = v, z = -1;
  for (int i = Dep[v]; i >= 0; --i) {
   Sub[x] += Dist[v][i], Cnt[x]++;
   if (z != -1)
    Sub2[z] += Dist[v][i], Cnt2[z]++;
   x = Pa[z = x];
 int64_t Query(int v) {
  int64_t res = 0;
  int x = v, z = -1;
  for (int i = Dep[v]; i >= 0; --i) {
   res += Sub[x] + 1LL * Cnt[x] * Dist[v][i];
   if (z != -1)
    res -= Sub2[z] + 1LL * Cnt2[z] * Dist[v][i];
   x = Pa[z = x];
  }
  return res;
}
};
```

#### 3.9 Lowbit Decomposition [760ac1]

```
int timer, chains;
vector<vector<int>> G;
vector<int> tl, tr, chain, head, dep, pa;
// chains : number of chain
// tl, tr[u] : subtree interval in the seq. of u
// head[i] : head of the chain i
// chian[u] : chain id of the chain u is on
void predfs(int u, int f) {
 dep[u] = dep[pa[u] = f] + 1;
 for (int v : G[u]) if (v != f) {
  predfs(v, u);
  if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
   chain[u] = chain[v];
 if (chain[u] == 0) chain[u] = ++chains;
void dfschain(int u, int f) {
```

```
tl[u] = timer++;
  if (head[chain[u]] == -1)
  head[chain[u]] = u;
  for (int v : G[u])
  if (v != f and chain[v] == chain[u])
    dfschain(v, u);
  for (int v : G[u])
  if (v != f and chain[v] != chain[u])
    dfschain(v, u);
 tr[u] = timer;
public:
LBD(\textbf{int}\ n)\ :\ timer(0),\ chains(0),\ G(n),\ tl(n),\ tr(n),
chain(n), head(n + 1, -1), dep(n), pa(n) \{\} void add_edge(int u, int v) \{
 G[u].push_back(v); G[v].push_back(u);
void decompose() { predfs(0, 0); dfschain(0, 0); }
PII get_subtree(int u) { return {tl[u], tr[u]}; }
vector<PII> get_path(int u, int v) {
 vector<PII> res;
 while (chain[u] != chain[v]) {
   if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
    swap(u, v);
   int s = head[chain[u]];
  res.emplace_back(tl[s], tl[u] + 1);
  u = pa[s];
 if (dep[u] < dep[v]) swap(u, v);</pre>
 res.emplace_back(tl[v], tl[u] + 1);
  return res;
};
3.10
      Virtual Tree [ad5cf5]
vector<pair<int, int>> build(vector<int> vs, int r) {
vector<pair<int, int>> res;
sort(vs.begin(), vs.end(), [](int i, int j) {
  return dfn[i] < dfn[j]; });</pre>
 vector<int> s = {r};
for (int v : vs) if (v != r) {
  if (int o = lca(v, s.back()); o != s.back()) {
  while (s.size() >= 2) {
   if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
    res.emplace_back(s[s.size() - 2], s.back());
    s.pop_back();
   if (s.back() != o) {
   res.emplace_back(o, s.back());
    s.back() = o;
  }
 }
  s.push_back(v);
for (size_t i = 1; i < s.size(); ++i)</pre>
  res.emplace_back(s[i - 1], s[i]);
 return res; // (x, y): x->y
3.11 Tree Hashing [707efa]
llu F(llu z) { // xorshift64star from iwiwi
z ^= z >> 12; z ^= z << 25; z ^= z >> 27;
return z * 2685821657736338717LL;
llu hsah(int u, int f) {
llu r = 127; // bigger?
for (int v : G[u]) if (v != f) r += F( hsah(v, u) );
return F(r):
} // test @ UOJ 763
3.12 Mo's Algorithm on Tree
dfs u:
push u
iterate subtree
Let P = LCA(u, v) with St(u) \le St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
      Count Cycles [c7e8f2]
```

// ord = sort by deg decreasing, rk[ord[i]] = i

// D[i] = edge point from rk small to rk big

```
for (int x : ord) { // c3
 for (int y : D[x]) vis[y] = 1;
 for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
 for (int y : D[x]) vis[y] = 0;
for (int x : ord) { // c4
 for (int y : D[x]) for (int z : adj[y])
  if (rk[z] > rk[x]) c4 += vis[z]++;
 for (int y : D[x]) for (int z : adj[y])
  if (rk[z] > rk[x]) --vis[z];
} // both are O(M*sqrt(M)), test @ 2022 CCPC guangzhou
3.14 MaximalClique [293730]
// contain a self loop u to u, than u won't in clique
template <size_t maxn> class MaxClique {
private:
 using bits = bitset<maxn>;
 bits popped, G[maxn], ans;
 size_t deg[maxn], deo[maxn], n;
 void sort_by_degree() {
  popped.reset();
  for (size_t i = 0; i < n; ++i)</pre>
   deg[i] = G[i].count();
  for (size_t i = 0; i < n; ++i) {
  size_t mi = maxn, id = 0;</pre>
   for (size_t j = 0; j < n; ++j)</pre>
    if (not popped[j] and deg[j] < mi)</pre>
     mi = deg[id = j];
   popped[deo[i] = id] = 1;
   for (size_t u = G[i]._Find_first(); u < n;</pre>
     u = G[i]._Find_next(u))
    --deg[u];
  }
 void BK(bits R, bits P, bits X) {
  if (R.count() + P.count() <= ans.count()) return;</pre>
  if (not P.count() and not X.count()) {
   if (R.count() > ans.count()) ans = R;
   return;
  }
  /* greedily chosse max degree as pivot
  bits cur = P | X; size_t pivot = 0, sz = 0;
  for ( size_t u = cur._Find_first() ;
   u < n ; u = cur.\_Find\_next(u))
   if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
  cur = P & ( ~G[ pivot ] );
  */ // or simply choose first
  bits cur = P & (~G[(P | X)._Find_first()]);
  for (size_t u = cur._Find_first(); u < n;</pre>
    u = cur._Find_next(u)) {
   if (R[u]) continue;
   R[u] = 1;
   BK(R, P \& G[u], X \& G[u]);
   R[u] = P[u] = 0, X[u] = 1;
 }
public:
 void init(size_t n_) {
  n = n_{;}
  for (size_t i = 0; i < n; ++i) G[i].reset();</pre>
  ans.reset();
 void add_edges(int u, bits S) { G[u] = S; }
 void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
 int solve() {
  sort_by_degree(); // or simply iota( deo... )
  for (size_t i = 0; i < n; ++i)</pre>
   deg[i] = G[i].count();
  bits pob, nob = 0; pob.set();
  for (size_t i = n; i < maxn; ++i) pob[i] = 0;</pre>
  for (size_t i = 0; i < n; ++i) {</pre>
   size_t v = deo[i];
   bits tmp;
   tmp[v] = 1;
   BK(tmp, pob \& G[v], nob \& G[v]);
   pob[v] = 0, nob[v] = 1;
  return static_cast<int>(ans.count());
};
```

#### 3.15 MaximumClique [aee5d8]

```
constexpr size_t kN = 150; using bits = bitset<kN>;
struct MaxClique {
bits G[kN], cs[kN];
 int ans, sol[kN], q, cur[kN], d[kN], n;
void init(int _n) {
 n = _n;
  for (int i = 0; i < n; ++i) G[i].reset();</pre>
void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
void pre_dfs(vector<int> &v, int i, bits mask) {
  if (i < 4) {
  for (int x : v) d[x] = (int)(G[x] \& mask).count();
  sort(all(v), [&](int x, int y) {
    return d[x] > d[y]; });
 vector<int> c(v.size());
  cs[1].reset(), cs[2].reset();
  int l = max(ans - q + 1, 1), r = 2, tp = 0, k;
  for (int p : v) {
   for (k = 1; (cs[k] & G[p]).any(); ++k);
   if (k >= r) cs[++r].reset();
   cs[k][p] = 1;
   if (k < l) v[tp++] = p;
 for (k = l; k < r; ++k)
  for (auto p = cs[k]._Find_first();</pre>
     p < kN; p = cs[k]._Find_next(p))
    v[tp] = (int)p, c[tp] = k, ++tp;
 dfs(v, c, i + 1, mask);
void dfs(vector<int> &v, vector<int> &c,
   int i, bits mask) {
  while (!v.empty()) {
   int p = v.back(); v.pop_back(); mask[p] = 0;
   if (q + c.back() <= ans) return;</pre>
   cur[q++] = p;
   vector<int> nr:
   for (int x : v) if (G[p][x]) nr.push_back(x);
   if (!nr.empty()) pre_dfs(nr, i, mask & G[p]);
   else if (q > ans) ans = q, copy_n(cur, q, sol);
   c.pop_back(); --q;
 }
int solve() {
 vector<int> v(n); iota(all(v), 0);
  ans = q = 0; pre_dfs(v, 0, bits(string(n, '1')));
  return ans; // sol[0 ~ ans-1]
} cliq; // test @ yosupo judge
```

# 3.16 Minimum Mean Cycle [e23bc0]

```
// WARNING: TYPE matters
struct Edge { int s, t; llf c; };
llf solve(vector<Edge> &e, int n) {
// O(VE), returns inf if no cycle, mmc otherwise
vector<VI> prv(n + 1, VI(n)), prve = prv;
vector<vector<llf>> d(n + 1, vector<llf>(n, inf));
d[0] = vector<llf>(n, 0);
 for (int i = 0; i < n; i++) {</pre>
  for (int j = 0; j < (int)e.size(); j++) {</pre>
   auto [s, t, c] = e[j];
if (d[i][s] < inf && d[i + 1][t] > d[i][s] + c) {
    d[i + 1][t] = d[i][s] + c;
    prv[i + 1][t] = s; prve[i + 1][t] = j;
  }
 llf mmc = inf; int st = -1;
 for (int i = 0; i < n; i++) {</pre>
  llf avg = -inf;
  for (int k = 0; k < n; k++) {
  if (d[n][i] < inf - eps)</pre>
    avg = max(avg, (d[n][i] - d[k][i]) / (n - k));
   else avg = inf;
  if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
if (st == -1) return inf;
vector<int> vst(n), eid, cycle, rho;
for (int i = n; !vst[st]; st = prv[i--][st]) {
  vst[st]++; eid.emplace_back(prve[i][st]);
  rho.emplace_back(st);
```

```
while (vst[st] != 2) {
  int v = rho.back(); rho.pop_back();
  cycle.emplace_back(v); vst[v]++;
}
reverse(all(eid)); eid.resize(cycle.size());
return mmc;
}
```

# 4 Flow & Matching 4.1 HopcroftKarp [930040]

```
struct HK {
 vector<int> l, r, a, p; int ans;
 HK(int n, int m, auto \&g) : l(n,-1), r(m,-1), ans(0) {
  for (bool match = true; match;) {
   match = false; a.assign(n, -1); p.assign(n, -1);
   queue<int> q;
   for (int i = 0; i < n; i++)</pre>
    if (l[i] == -1) q.push(a[i] = p[i] = i);
   // bitset<maxn> nvis, t; nvis.set();
   while (!q.empty()) {
    int z, x = q.front(); q.pop();
    if (l[a[x]] != -1) continue;
    for (int y : g[x]) { // or iterate t = g[x]&nvis
     // nvis.reset(y);
     if (r[y] == -1) {
      for (z = y; z != -1;)
       r[z] = x, swap(l[x], z), x = p[x];
      match = true; ++ans; break;
     } else if (p[r[y]] == -1)
      q.push(z = r[y]), p[z] = x, a[z] = a[x];
 }
};
```

# 4.2 Dijkstra Cost Flow [06d5f2]

```
template <typename F, typename C> class MCMF {
 static constexpr F INF_F = numeric_limits<F>::max();
 static constexpr C INF_C = numeric_limits<C>::max();
 struct E {
  int to, r; F f; C c;
 E(int a, int b, F x, C y)
   : to(a), r(b), f(x), c(y) {}
 vector<vector<E>> g; vector<pair<int, int>> f;
 vector<F> up; vector<C> d, h;
 optional<pair<F, C>> step(int S, int T) {
 priority_queue<pair<C, int>> q;
  q.emplace(d[S] = 0, S), up[S] = INF_F;
  while (not q.empty()) {
   auto [l, u] = q.top(); q.pop();
   if (up[u] == 0 or l != -d[u]) continue;
   for (int i = 0; i < int(g[u].size()); ++i) {</pre>
    auto e = g[u][i]; int v = e.to;
    auto nd = d[u] + e.c + h[u] - h[v];
    if (e.f <= 0 or d[v] <= nd) continue;</pre>
    f[v] = \{u, i\}; up[v] = min(up[u], e.f);
    q.emplace(-(d[v] = nd), v);
  if (d[T] == INF_C) return nullopt;
  for (size_t i = 0; i < d.size(); i++) h[i]+=d[i];</pre>
  for (int i = T; i != S; i = f[i].first) {
  auto &eg = g[f[i].first][f[i].second];
   eg.f -= up[T]; g[eg.to][eg.r].f += up[T];
  return pair{up[T], h[T]};
public:
MCMF(int n) : g(n), f(n), up(n), d(n, INF_C) {}
 void add_edge(int s, int t, F c, C w) {
 g[s].emplace_back(t, int(g[t].size()), c, w);
 g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
 pair<F, C> solve(int a, int b) {
 h.assign(g.size(), 0);
  F c = 0; C w = 0;
  while (auto r = step(a, b)) {
   c += r->first, w += r->first * r->second;
   fill(d.begin(), d.end(), INF_C);
```

```
return {c, w};
}
};
```

#### Dinic [32c53e]

```
template <typename Cap = int64_t> class Dinic {
private:
 struct E { int to, rev; Cap cap; }; int n, st, ed;
 vector<vector<E>> G; vector<size_t> lv, idx;
 bool BFS(int k) {
  lv.assign(n, 0); idx.assign(n, 0);
  queue<int> bfs; bfs.push(st); lv[st] = 1;
  while (not bfs.empty() and not lv[ed]) {
   int u = bfs.front(); bfs.pop();
for (auto e: G[u]) if (e.cap >> k and !lv[e.to])
    bfs.push(e.to), lv[e.to] = lv[u] + 1;
  return lv[ed];
 Cap DFS(int u, Cap f = numeric_limits<Cap>::max()) {
  if (u == ed) return f;
  Cap ret = 0;
  for (auto &i = idx[u]; i < G[u].size(); ++i) {</pre>
   auto &[to, rev, cap] = G[u][i];
if (cap <= 0 or lv[to] != lv[u] + 1) continue;</pre>
   Cap nf = DFS(to, min(f, cap));
   ret += nf; cap -= nf; f -= nf;
   G[to][rev].cap += nf;
   if (f == 0) return ret;
  if (ret == 0) lv[u] = 0;
  return ret;
public:
 void init(int n_) { G.assign(n = n_, vector<E>()); }
 void add_edge(int u, int v, Cap c) {
  G[u].push_back({v, int(G[v].size()), c});
G[v].push_back({u, int(G[u].size())-1, 0});
 Cap max_flow(int st_, int ed_) {
  st = st_, ed = ed_; Cap ret = 0;

for (int i = 63; i >= 0; --i)
   while (BFS(i)) ret += DFS(st);
  return ret;
}; // test @ luogu P3376
```

### 4.4 Flow Models

- · Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x, y, l, u), connect  $x \to y$  with capacity u l.
  - For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, con- $\operatorname{nect} v \to T$  with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer. Also, fis a mincost valid flow.
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e\,+\,f_e$  , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching  ${\cal M}$  on bi-
  - 1. Redirect every edge:  $y \to x$  if  $(x, y) \in M$ ,  $x \to y$  otherwise.
  - 2. DFS from unmatched vertices in  $\hat{X}$ .
  - 3.  $x \in X$  is chosen iff x is unvisited. 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source S and sink T
  - 2. For each edge (x, y, c), connect  $x \to y$  with (cost, cap) = (c, 1) if c>0, otherwise connect y o x with (cost, cap)=(-c,1)
  - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1</li>
    4. For each vertex v with d(v) > 0, connect S → v with (cost, cap) =

  - 5. For each vertex v with d(v) < 0, connect v  $\rightarrow$  T with (cost, cap) = (0, -d(v)) 6. Flow from S to T , the answer is the cost of the flow C+K
- · Maximum density induced subgraph

- 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
- 2. Construct a max flow model, let K be the sum of all weights
- 3. Connect source  $s \to v$ ,  $v \in G$  with capacity K
- 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
- 5. For  $v \in \mathit{G}$ , connect it with sink  $v \to t$  with capacity K + 2T - $\left(\sum_{e \in E(v)} w(e)\right) - 2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with
  - weight w(u,v). 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to  $\emph{\emph{v}}$  .
  - 3. Find the minimum weight perfect matching on G'.
- · Submodular functions minimization
  - For a function  $f: 2^V \to \mathbb{R}$ , f is a submodular function iff

```
* \forall S, T \subseteq V, f(S) + f(T) \ge f(S \cup T) + f(S \cap T), or
```

- $* \ \forall X \subseteq Y \subseteq V, x \notin Y, f(X \cup \{x\}) f(X) \geq f(Y \cup \{x\}) f(Y).$
- To minimize  $\sum_i \theta_i(x_i)$  $\sum_{i < j} \phi_{ij}(x_i, x_j)$  $\sum_{i < j < k} \psi_{ijk}(x_i, x_j, x_k)$
- If  $\theta_i(1) \geq \theta_i(0)$ , add edge  $(S, i, \theta_i(1) \theta_i(0))$  and  $\theta_i(0)$  to answer; otherwise, (i, T,  $\theta_i(0) - \theta_i(1)$ ) and  $\theta_i(1)$ .
- Add edges  $(i, j, \phi_{ij}(0, 1) + \phi_{ij}(1, 0) \phi_{ij}(0, 0) \phi_{ij}(1, 1)).$
- Denote  $x_{ijk}$  as helper nodes. Let  $P = \psi_{ijk}(0,0,0) + \psi_{ijk}(0,1,1) + \psi_{ijk}(1,0,1) + \psi_{ijk}(1,1,0) \psi_{ijk}(0,0,1) \psi_{ijk}(0,1,0) \psi_{ijk}(1,0,0) \psi_{ijk}(1,1,1)$ . Add -P to answer. If  $P \geq 0$ , add edges  $(i,x_{ijk},P)$ ,  $(j,x_{ijk},P)$ ,  $(k,x_{ijk},P)$ ,  $(x_{ijk},T,P)$ ; otherwise  $(x_{ijk},i,-P)$ ,  $(x_{ijk},j,-P)$ ,  $(x_{ijk},k,-P)$ ,  $(S,x_{ijk},-P)$ .
- The minimum cut of this graph will be the the minimum value of the function above.

# 4.5 General Graph Matching [5f2293]

struct Matching {

```
queue<int> q; int ans, n;
vector<int> fa, s, v, pre, match;
 int Find(int u) {
  return u == fa[u] ? u : fa[u] = Find(fa[u]); }
 int LCA(int x, int y) {
  static int tk = 0; tk++; x = Find(x); y = Find(y);
  for (;; swap(x, y)) if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
  }
 void Blossom(int x, int y, int l) {
  for (; Find(x) != l; x = pre[y]) {
   pre[x] = y, y = match[x];
   if (s[y] == 1) q.push(y), s[y] = 0;
   for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
  }
 bool Bfs(auto &&g, int r) {
  iota(all(fa), 0); ranges::fill(s, -1);
  q = queue<int>(); q.push(r); s[r] = 0;
  for (; !q.empty(); q.pop()) {
   for (int x = q.front(); int u : g[x])
    if (s[u] == -1) {
     if (pre[u] = x, s[u] = 1, match[u] == n) {
      for (int a = u, b = x, last;
        b != n; a = last, b = pre[a])
       last = match[b], match[b] = a, match[a] = b;
      return true;
     q.push(match[u]); s[match[u]] = 0;
    } else if (!s[u] && Find(u) != Find(x)) {
     int l = LCA(u, x);
     Blossom(x, u, l); Blossom(u, x, l);
    }
  return false;
 Matching(auto &&g) : ans(0), n(int(g.size())),
 fa(n+1), s(n+1), v(n+1), pre(n+1, n), match(n+1, n) {
  for (int x = 0; x < n; ++x)
   if (match[x] == n) ans += Bfs(g, x);
 } // match[x] == n means not matched
}; // test @ yosupo judge
```

# 4.6 Global Min-Cut [1f0306]

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
pair<int, int> phase(int n) {
 memset(v, false, sizeof(v));
 memset(g, 0, sizeof(g));
 int s = -1, t = -1;
 while (true) {
  int c = -1;
  for (int i = 0; i < n; ++i) {</pre>
   if (del[i] || v[i]) continue;
   if (c == -1 || g[i] > g[c]) c = i;
  if (c == -1) break;
  v[s = t, t = c] = true;
  for (int i = 0; i < n; ++i) {
   if (del[i] || v[i]) continue;
   g[i] += w[c][i];
 return make_pair(s, t);
int mincut(int n) {
 int cut = 1e9;
 memset(del, false, sizeof(del));
 for (int i = 0; i < n - 1; ++i) {</pre>
 int s, t; tie(s, t) = phase(n);
del[t] = true; cut = min(cut, g[t]);
  for (int j = 0; j < n; ++j) {</pre>
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
  }
 }
 return cut;
```

# 4.7 GomoryHu Tree [7473bb]

```
vector<tuple<int, int, int>> GomoryHu(int n){
vector<tuple<int, int, int>> rt;
vector<int> g(n);
for (int i = 1; i < n; ++i) {</pre>
 int t = g[i];
 flow.reset(); // clear flows on all edge
  rt.emplace_back(i, t, flow.max_flow(i, t));
 flow.walk(i); // bfs points that connected to i (use
 edges with .cap > 0)
for (int j = i + 1; j < n; ++j)</pre>
  if (g[j]==t&&flow.connect(j)) // check if i can
    reach j
   g[j] = i;
}
return rt;
```

#### 4.8 Kuhn Munkres [2c09ed]

```
struct KM { // maximize, test @ UOJ 80
int n, l, r; lld ans; // fl and fr are the match
vector<lld> hl, hr; vector<int> fl, fr, pre, q;
void bfs(const auto &w, int s) {
 vector<int> vl(n), vr(n); vector<lld> slk(n, INF);
 l = r = 0; vr[q[r++] = s] = true;
 const auto check = [&](int x) -> bool {
  if (vl[x] || slk[x] > 0) return true;
  vl[x] = true; slk[x] = INF;
   if (fl[x] != -1) return vr[q[r++] = fl[x]] = true;
  while (x != -1) swap(x, fr[fl[x] = pre[x]]);
  return false;
 }:
 while (true) {
   for (int x = 0, y = q[l++]; x < n; ++x) if (!vl[x])
    if (chmin(slk[x], hl[x] + hr[y] - w[x][y]))
      if (pre[x] = y, !check(x)) return;
  lld d = ranges::min(slk);
   for (int x = 0; x < n; ++x)
   vl[x] ? hl[x] += d : slk[x] -= d;
   for (int x = 0; x < n; ++x) if (vr[x]) hr[x] -= d;
   for (int x = 0; x < n; ++x) if (!check(x)) return;
```

```
KM(int n_, const auto &w) : n(n_), ans(0),
  hl(n), hr(n), fl(n, -1), fr(fl), pre(n), q(n) {
  for (int i = 0; i < n; ++i) hl[i]=ranges::max(w[i]);
for (int i = 0; i < n; ++i) bfs(w, i);</pre>
  for (int i = 0; i < n; ++i) ans += w[i][fl[i]];</pre>
};
```

```
4.9 Minimum Cost Circulation [0f0e85]
int vis[N], visc, fa[N], fae[N], head[N], mlc = 1;
struct ep {
 int to, next;
ll flow, cost;
} e[M << 1];
void adde(int u, int v, ll fl, int cs) {
 e[++mlc] = {v, head[u], fl, cs};
head[u] = mlc;
 e[++mlc] = {u, head[v], 0, -cs};
head[v] = mlc;
void dfs(int u) {
 vis[u] = 1;
 for (int i = head[u], v; i; i = e[i].next)
 if (!vis[v = e[i].to] and e[i].flow)
   fa[v] = u, fae[v] = i, dfs(v);
ll phi(int x) {
static ll pi[N];
 if (x == -1) return 0;
 if (vis[x] == visc) return pi[x];
 return vis[x] = visc, pi[x] = phi(fa[x]) - e[fae[x]].
void pushflow(int x, ll &cost) {
 int v = e[x ^ 1].to, u = e[x].to;
 ++visc:
 while (v != -1) vis[v] = visc, v = fa[v];
 while (u != -1 && vis[u] != visc)
 vis[u] = visc, u = fa[u];
 vector<int> cyc;
 int e2 = 0, pa = 2;
 ll f = e[x].flow;
 for (int i = e[x ^ 1].to; i != u; i = fa[i]) {
  cyc.push_back(fae[i]);
  if (e[fae[i]].flow < f)</pre>
   f = e[fae[e2 = i] ^ (pa = 0)].flow;
 for (int i = e[x].to; i != u; i = fa[i]) {
 cyc.push_back(fae[i] ^ 1);
  if (e[fae[i] ^ 1].flow < f)</pre>
   f = e[fae[e2 = i] ^ (pa = 1)].flow;
 cyc.push_back(x);
 for (int cyc_i : cyc) {
  e[cyc_i].flow -= f, e[cyc_i ^ 1].flow += f;
  cost += 1ll * f * e[cyc_i].cost;
 if (pa == 2) return;
 int le = x ^ pa, l = e[le].to, o = e[le ^ 1].to;
 while (l != e2) {
 vis[o] = 0;
  swap(le ^= 1, fae[o]), swap(l, fa[o]), swap(l, o);
}
ll simplex() { // 1-based
 ll cost = 0;
 memset(fa, -1, sizeof(fa)), dfs(1);
 vis[1] = visc = 2, fa[1] = -1;
 for (int i = 2, pre = -1; i != pre; i = (i == mlc ? 2
    : i + 1))
  if (e[i].flow and e[i].cost < phi(e[i ^ 1].to) - phi(</pre>
    e[i].to))
   pushflow(pre = i, cost);
 return cost;
}
```

# 4.10 Minimum Cost Max Flow [6d1b01]

```
template <typename F, typename C> class MCMF {
 static constexpr F INF_F = numeric_limits<F>::max();
 static constexpr C INF_C = numeric_limits<C>::max();
 struct E {
```

```
int to, r;
  F f; C c;
  E() {}
  E(int a, int b, F x, C y)
: to(a), r(b), f(x), c(y) {}
 };
 vector<vector<E>> g;
 vector<pair<int, int>> f;
 vector<bool> inq;
 vector<F> up; vector<C> d;
 optional<pair<F, C>> step(int S, int T) {
  queue<int> q;
  for (q.push(S), d[S] = 0, up[S] = INF_F;
    not q.empty(); q.pop()) {
   int u = q.front(); inq[u] = false;
   if (up[u] == 0) continue;
for (int i = 0; i < int(g[u].size()); ++i) {</pre>
    auto e = g[u][i]; int v = e.to;
    if (e.f <= 0 or d[v] <= d[u] + e.c)
     continue;
    d[v] = d[u] + e.c; f[v] = {u, i};
    up[v] = min(up[u], e.f);
    if (not inq[v]) q.push(v);
    inq[v] = true;
   }
  if (d[T] == INF_C) return nullopt;
  for (int i = T; i != S; i = f[i].first) {
   auto &eg = g[f[i].first][f[i].second];
   eg.f -= up[T];
   g[eg.to][eg.r].f += up[T];
  return pair{up[T], d[T]};
public:
 MCMF(int n) : g(n), f(n), inq(n), up(n), d(n, INF_C) {}
 void add_edge(int s, int t, F c, C w) {
  g[s].emplace_back(t, int(g[t].size()), c, w);
g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
 pair<F, C> solve(int a, int b) {
  F c = 0; C w = 0;
  while (auto r = step(a, b)) {
   c += r->first, w += r->first * r->second;
   fill(inq.begin(), inq.end(), false);
   fill(d.begin(), d.end(), INF_C);
  return {c, w};
};
```

## 4.11 Weighted Matching [94ca35]

```
#define pb emplace_back
#define rep(i, l, r) for (int i=(l); i<=(r); ++i)
struct WeightGraph { // 1-based
  static const int inf = INT_MAX;
struct edge { int u, v, w; }; int n, nx;
vector<int> lab; vector<vector<edge>> g;
vector<int> slack, match, st, pa, S, vis;
vector<vector<int>>> flo, flo_from; queue<int> q;
WeightGraph(int n_) : n(n_-), nx(n * 2), lab(nx + 1), g(nx + 1, vector < edge > (nx + 1), slack(nx + 1),
  flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
  match = st = pa = S = vis = slack;
  rep(u, 1, n) rep(v, 1, n) g[u][v] = {u, v, 0};
int ED(edge e) {
return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
void update_slack(int u, int x, int &s) {
  if (!s || ED(g[u][x]) < ED(g[s][x])) s = u; }</pre>
void set_slack(int x) {
  slack[x] = 0;
  for (int u = 1; u <= n; ++u)
   if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
    update_slack(u, x, slack[x]);
void q_push(int x) {
  if (x \le n) q.push(x);
  else for (int y : flo[x]) q_push(y);
void set_st(int x, int b) {
 st[x] = b;
```

```
if (x > n) for (int y : flo[x]) set_st(y, b);
vector<int> split_flo(auto &f, int xr) {
 auto it = find(all(f), xr);
 if (auto pr = it - f.begin(); pr % 2 == 1)
 reverse(1 + all(f)), it = f.end() - pr;
 auto res = vector(f.begin(), it);
 return f.erase(f.begin(), it), res;
void set_match(int u, int v) {
 match[u] = g[u][v].v;
 if (u <= n) return;</pre>
 int xr = flo_from[u][g[u][v].u];
 auto &f = flo[u], z = split_flo(f, xr);
 rep(i, 0, int(z.size())-1) set_match(z[i], z[i ^ 1]);
 set_match(xr, v); f.insert(f.end(), all(z));
void augment(int u, int v) {
 for (;;) {
  int xnv = st[match[u]]; set_match(u, v);
  if (!xnv) return;
  set_match(xnv, st[pa[xnv]]);
  u = st[pa[xnv]], v = xnv;
int lca(int u, int v) {
  static int t = 0; ++t;
 for (++t; u || v; swap(u, v)) if (u) {
  if (vis[u] == t) return u;
  vis[u] = t; u = st[match[u]];
  if (u) u = st[pa[u]];
 return 0;
void add_blossom(int u, int o, int v) {
 int b = int(find(n + 1 + all(st), 0) - begin(st));
 lab[b] = 0, S[b] = 0; match[b] = match[o];
 vector<int> f = {o};
 for (int x = u, y; x != o; x = st[pa[y]])
f.pb(x), f.pb(y = st[match[x]]), q_push(y);
 reverse(1 + all(f));
 for (int x = v, y; x != o; x = st[pa[y]])
f.pb(x), f.pb(y = st[match[x]]), q_push(y);
 flo[b] = f; set_st(b, b);
 for (int x = 1; x <= nx; ++x)
  g[b][x].w = g[x][b].w = 0;
 for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;</pre>
 for (int xs : flo[b]) {
  for (int x = 1; x <= nx; ++x)
   if (g[b][x].w == 0 \mid \mid ED(g[xs][x]) < ED(g[b][x]))
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
  for (int x = 1; x \le n; ++x)
   if (flo_from[xs][x]) flo_from[b][x] = xs;
 }
 set_slack(b);
void expand_blossom(int b) {
 for (int x : flo[b]) set_st(x, x);
 int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
 for (int x : split_flo(flo[b], xr)) {
  if (xs == -1) { xs = x; continue; }
  pa[xs] = g[x][xs].u; S[xs] = 1, S[x] = 0;
  slack[xs] = 0; set_slack(x); q_push(x); xs = -1;
 for (int x : flo[b])
  if (x == xr) S[x] = 1, pa[x] = pa[b];
  else S[x] = -1, set_slack(x);
 st[b] = 0;
bool on_found_edge(const edge &e) {
 if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
  int nu = st[match[v]]; pa[v] = e.u; S[v] = 1;
  slack[v] = slack[nu] = 0; S[nu] = 0; q_push(nu);
 } else if (S[v] == 0) {
  if (int o = lca(u, v)) add_blossom(u, o, v);
  else return augment(u, v), augment(v, u), true;
 return false;
bool matching() {
 ranges::fill(S, -1); ranges::fill(slack, 0);
 q = queue<int>();
```

```
for (int x = 1; x <= nx; ++x)
   if (st[x] == x && !match[x])
    pa[x] = 0, S[x] = 0, q_push(x);
  if (q.empty()) return false;
  for (;;) {
   while (q.size()) {
    int u = q.front(); q.pop();
    if (S[st[u]] == 1) continue;
    for (int v = 1; v <= n; ++v)</pre>
     if (g[u][v].w > 0 && st[u] != st[v]) {
      if (ED(g[u][v]) != 0)
       update_slack(u, st[v], slack[st[v]]);
      else if (on_found_edge(g[u][v])) return true;
   int d = inf;
   for (int b = n + 1; b <= nx; ++b)
   if (st[b] == b && S[b] == 1)
     d = min(d, lab[b] / 2);
   for (int x = 1; x <= nx; ++x)</pre>
    if (int s = slack[x]; st[x] == x && s && S[x] <= 0)</pre>
     d = min(d, ED(g[s][x]) / (S[x] + 2));
   for (int u = 1; u <= n; ++u)
if (S[st[u]] == 1) lab[u] += d;</pre>
    else if (S[st[u]] == 0) {
     if (lab[u] <= d) return false;</pre>
     lab[u] -= d;
   rep(b, n + 1, nx) if (st[b] == b \&\& S[b] >= 0)
    lab[b] += d * (2 - 4 * S[b]);
   for (int x = 1; x \le nx; ++x)
    if (int s = slack[x]; st[x] == x &&
      s \&\& st[s] != x \&\& ED(g[s][x]) == 0)
     if (on_found_edge(g[s][x])) return true;
   for (int b = n + 1; b <= nx; ++b)</pre>
    if (st[b] == b && S[b] == 1 && lab[b] == 0)
     expand_blossom(b);
  return false;
pair<lld, int> solve() {
  ranges::fill(match, 0);
  rep(u, 0, n) st[u] = u, flo[u].clear();
  int w_max = 0;
  rep(u, 1, n) rep(v, 1, n) {
  flo_from[u][v] = (u == v ? u : 0);
   w_max = max(w_max, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
  int n_matches = 0; lld tot_weight = 0;
  while (matching()) ++n_matches;
  rep(u, 1, n) if (match[u] \&\& match[u] < u)
  tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
void set_edge(int u, int v, int w) {
  g[u][v].w = g[v][u].w = w; }
5
     Math
5.1 Common Bounds
```

$$\begin{split} p(0) &= 1, \ p(n) = \sum_{k \in \mathbb{Z} \backslash \{0\}} (-1)^{k+1} p(n-k(3k-1)/2) \approx 0.145/n \cdot \exp(2.56\sqrt{n}) \\ &\frac{n}{\max_{i \leq n} (d(i))} \frac{100 \text{ le3 le6 le9 le12 le15 le18}}{12 \text{ 32 240 1344 6720 26880 103680}} \end{split}$$

#### 5.2 Stirling Number

First Kind  $S_1(n,k)$  counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1) \dots (x+n-1) = \sum_{k=0}^{n} S_1(n,k) x^k$$

$$g(x) = x(x+1) \dots (x+n-1) = \sum_{k=0}^{n} a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^{n} \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^{k} ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

#### Second Kind

 $S_2(n,k)$  counts the number of ways to partition a set of n elements into knonempty sets.  $S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$ 

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

# 5.3 ax+by=gcd [d0cbdd]

```
// ax+ny = 1, ax+ny == ax == 1 \pmod{n}
void exgcd(lld x, lld y, lld &g, lld &a, lld &b) {
 if (y == 0) g = x, a = 1, b = 0;
 else exgcd(y, x \% y, g, b, a), b -= (x / y) * a;
```

# 5.4 Chinese Remainder [d69e74]

```
// please ensure r_i \in [0, m_i)
bool crt(lld &m1, lld &r1, lld m2, lld r2) {
  if (m2 > m1) swap(m1, m2), swap(r1, r2);
  lld g, a, b; exgcd(m1, m2, g, a, b);
  if ((r2 - r1) % g != 0) return false;
  m2 /= g; lld D = (r2 - r1) / g % m2 * a % m2;
r1 += (D < 0 ? D + m2 : D) * m1; m1 *= m2;
  assert (r1 >= 0 \&\& r1 < m1);
  return true;
```

# 5.5 DiscreteLog [86e463]

```
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
 // x^? \setminus equiv y \pmod{M}
 Int t = 1, c = 0, g = 1;
for (Int M_ = M; M_ > 0; M_ >>= 1) g = g * x % M;
for (g = gcd(g, M); t % g != 0; ++c) {
   if (t == y) return c;
  t = t * x % M;
 if (y % g != 0) return -1;
 t /= g, y /= g, M /= g;
 Int h = 0, gs = 1;
 for (; h * h < M; ++h) gs = gs * x % M;
 unordered_map<Int, Int> bs;
 for (Int s = 0; s < h; bs[y] = ++s) y = y * x % M;
 for (Int s = 0; s < M; s += h) {</pre>
  t = t * gs % M;
  if (bs.count(t)) return c + s + h - bs[t];
 return -1;
```

#### 5.6 Quadratic Residue [leabad]

```
int get_root(int n, int P) { // ensure 0 <= n < p</pre>
 if (P == 2 or n == 0) return n;
 auto check = [&](int x) {
  return modpow(x, (P - 1) / 2, P); };
if (check(n) != 1) return -1;
mt19937 rnd(7122); lld z = 1, w;
 while (check(w = (z * z - n + P) % P) != P - 1)
 z = rnd() \% P;
 const auto M = [P, w](auto &u, auto &v) {
  auto [a, b] = u; auto [c, d] = v;
  return make_pair((a * c + b * d % P * w) % P,
    (a * d + b * c) \% P);
 pair<lld, lld> r(1, 0), e(z, 1);
 for (int w = (P + 1) / 2; w; w >>= 1, e = M(e, e))
  if (w & 1) r = M(r, e);
 return r.first; // sqrt(n) mod P where P is prime
```

#### 5.7 Extended Euler

$$a^b \equiv \begin{cases} a^{(b \mod \varphi(m)) + \varphi(m)} & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m$$

### 5.8 Extended FloorSum

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

```
\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}
```

#### 5.9 Extended Euclidean [e09892]

```
template <typename T>
auto euclid(lld a, lld b, lld c, lld n, T U, T R) {
b %= c;
if (a >= c)
 return euclid(a % c, b, c, n, U, mpow(U, a / c) * R);
lld m = (i128(a) * n + b) / c;
if (!m) return mpow(R, n);
return mpow(R, (c - b - 1) / a) * U *
       euclid(c, c - b - 1, a, m - 1, R, U) \star
       mpow(R, n - (i128(c) * m - b - 1) / a);
// time complexity is O(log max(a, c))
// 給定二維座標系上的一次函數 $y = (ax + b) / c$
// 維護一個矩陣 $A = I$, 考慮 $x \in [0, n)$
// 每次向右穿過網格的垂直線時,乘上一個矩陣 $R$
// 每次向上穿過網格的水平線時,乘上一個矩陣 $U$
// 若剛好經過一個整點,那麼先乘 $U$ 再乘 $R$
```

#### 5.10 FloorSum [fb5917]

```
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
// @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
llu ans = 0;
while (true) {
 if (a >= m) ans += n*(n-1)/2 * (a/m), a %= m;
  if (b >= m) ans += n * (b/m), b %= m;
 if (llu y_max = a * n + b; y_max >= m) {
  n = (llu)(y_max / m), b = (llu)(y_max % m);
  swap(m, a);
 } else break;
return ans;
lld floor_sum(lld n, lld m, lld a, lld b) {
llu ans = 0;
if (a < 0) {
 llu a2 = (a \% m + m), d = (a2 - a) / m;
 ans -= 1ULL * n * (n - 1) / 2 * d; a = a2;
if (b < 0) {
 llu b2 = (b \% m + m), d = (b2 - b) / m;
 ans -= 1ULL * n * d; b = b2;
return ans + floor_sum_unsigned(n, m, a, b);
```

#### 5.11 ModMin [253e4d]

```
// min{k | l <= ((ak) mod m) <= r}
optional<llu> mod_min(u32 a, u32 m, u32 l, u32 r) {
   if (a == 0) return l ? nullopt : 0;
   if (auto k = llu(l + a - 1) / a; k * a <= r)
      return k;
   auto b = m / a, c = m % a;
   if (auto y = mod_min(c, a, a - r % a, a - l % a))
    return (l + *y * c + a - 1) / a + *y * b;
   return nullopt;
}</pre>
```

#### **5.12** FWT [c5167a]

```
/* or convolution:
    * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
    * and convolution:
    * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
void fwt(int x[], int N, bool inv = false) {
    for (int d = 1; d < N; d <<= 1)
    for (int s = 0; s < N; s += d * 2)
    for (int i = s; i < s + d; i++) {
        int j = i + d, ta = x[i], tb = x[j];
    }
}</pre>
```

```
x[i] = modadd(ta, tb);
x[j] = modsub(ta, tb);
}
if (inv) {
  const int invn = modinv(N);
  for (int i = 0; i < N; i++)
   x[i] = modmul(x[i], invn);
}</pre>
```

# 5.13 Packed FFT [321552]

```
int round2k(size_t n) {
 int sz = 1; while (sz < int(n)) sz *= 2; return sz; }</pre>
VL convolution(const VI &a, const VI &b) {
 const int sz = round2k(a.size() + b.size() - 1);
 // Should be able to handle N <= 10^5, C <= 10^4
 vector<P> v(sz);
 for (size_t i = 0; i < a.size(); i++) v[i].RE(a[i]);</pre>
 for (size_t i = 0; i < b.size(); i++) v[i].IM(b[i]);</pre>
 fft(v.data(), sz, /*inv=*/false);
auto rev = v; reverse(1 + all(rev));
 for (int i = 0; i < sz; i++) {</pre>
  P A = (v[i] + conj(rev[i])) / P(2, 0);
  P B = (v[i] - conj(rev[i])) / P(0, 2);
  v[i] = A * B;
 VL c(sz); fft(v.data(), sz, /*inv=*/true);
 for (int i = 0; i < sz; ++i) c[i] = roundl(RE(v[i]));</pre>
 return c;
VI convolution_mod(const VI &a, const VI &b) {
 const int sz = round2k(a.size() + b.size() - 1);
 vector<P> fa(sz), fb(sz);
 for (size_t i = 0; i < a.size(); ++i)</pre>
  fa[i] = P(a[i] & ((1 << 15) - 1), a[i] >> 15);
 for (size_t i = 0; i < b.size(); ++i)
fb[i] = P(b[i] & ((1 << 15) - 1), b[i] >> 15);
 fft(fa.data(), sz); fft(fb.data(), sz);
 auto rfa = fa; reverse(1 + all(rfa));
 for (int i = 0; i < sz; ++i) fa[i] *= fb[i];</pre>
 for (int i = 0; i < sz; ++i) fb[i] *= conj(rfa[i]);</pre>
 fft(fa.data(), sz, true); fft(fb.data(), sz, true);
 vector<int> res(sz);
 for (int i = 0; i < sz; ++i) {</pre>
  lld A = (lld)roundl(RE((fa[i] + fb[i]) / P(2, 0)));
lld C = (lld)roundl(IM((fa[i] - fb[i]) / P(0, 2)));
  lld B = (lld)roundl(IM(fa[i])); B %= p; C %= p;
  res[i] = (A + (B << 15) + (C << 30)) % p;
 return res;
} // test @ yosupo judge with long double
```

#### 5.14 CRT for arbitrary mod [e4dde7]

```
const int mod = 10000000007;
const int M1 = 985661441; // G = 3 for M1, M2, M3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(lld A, lld B, lld C) {
    static_assert (M1 < M2 && M2 < M3);
    constexpr lld r12 = modpow(M1, M2-2, M2);
    constexpr lld r13 = modpow(M1, M3-2, M3);
    constexpr lld r23 = modpow(M2, M3-2, M3);
    constexpr lld M1M2 = 1LL * M1 * M2 % mod;
    B = (B - A + M2) * r12 % M2;
    C = (C - A + M3) * r13 % M3;
    C = (C - B + M3) * r23 % M3;
    return (A + B * M1 + C * M1M2) % mod;
}</pre>
```

#### **5.15** NTT / FFT [03190d]

```
template <int mod, int G, int maxn> struct NTT {
    static_assert (maxn == (maxn & -maxn));
    int roots[maxn];
    NTT () {
        int r = modpow(G, (mod - 1) / maxn);
        for (int i = maxn >> 1; i; i >>= 1) {
            roots[i] = 1;
            for (int j = 1; j < i; j++)
            roots[i + j] = modmul(roots[i + j - 1], r);
            r = modmul(r, r);
            // for (int j = 0; j < i; j++) // FFT (tested)
            // roots[i+j] = polar<llf>(1, PI * j / i);
```

Poly Sqrt(const Poly &v) { // need: QuadraticResidue

assert(!v.empty() && v[0] != 0);

```
}
                                                               const int r = get_root(v[0]); assert(r != -1);
                                                               return Newton(v, [r](int x) { return r; },
 // n must be 2^k, and 0 \le F[i] \le mod
                                                                [](Poly &X, Poly &A, int sz) {
 void operator()(int F[], int n, bool inv = false) {
for (int i = 0, j = 0; i < n; i++) {</pre>
                                                                 auto Y = X; Y.resize(sz / 2);
                                                                 auto B = Mul(A, Inv(Y), sz);
   if (i < j) swap(F[i], F[j]);</pre>
                                                                 for (int i = 0, inv2 = mod / 2 + 1; i < sz; i++)</pre>
   for (int k = n>>1; (j^=k) < k; k>>=1);
                                                                  X[i] = modmul(inv2, modadd(X[i], B[i])); });
  for (int s = 1; s < n; s *= 2) {
                                                              Poly Mul(auto &&a, auto &&b) {
   for (int i = 0; i < n; i += s * 2) {
                                                               const int n = a.size() + b.size() - 1;
    for (int j = 0; j < s; j++) {
                                                               auto R = Mul(a, b, bit_ceil(n));
     int a = F[i+j], b = modmul(F[i+j+s], roots[s+j]);
                                                               return R.resize(n), R;
     F[i+j] = modadd(a, b); // a + b
     F[i+j+s] = modsub(a, b); // a - b
                                                              Poly MulT(Poly a, Poly b, int k) {
                                                               assert(b.size()); reverse(all(b)); auto R = Mul(a, b);
                                                               R = vector(R.begin() + b.size() - 1, R.end());
   }
                                                               return R.resize(k), R;
  if (inv) {
   int iv = modinv(n);
                                                              Poly Eval(const Poly &f, const Poly &x) {
   for (int i = 0; i < n; i++) F[i] = modmul(F[i], iv);</pre>
                                                               if (f.empty()) return vector(x.size(), 0);
   reverse(F + 1, F + n);
                                                               const int n = int(max(x.size(), f.size()));
                                                               auto q = vector(n * 2, Poly(2, 1)); Poly ans(n);
                                                               fi(0, x.size()) q[i + n][1] = modsub(0, x[i]);
};
                                                               for (int i = n - 1; i > 0; i--)
                                                                q[i] = Mul(q[i << 1], q[i << 1 | 1]);
5.16
      Formal Power Series [2bc0d3]
                                                               q[1] = MulT(f, Inv(q[1]), n);
#define fi(l, r) for (size_t i = (l); i < (r); ++i)
                                                               for (int i = 1; i < n; i++) {</pre>
using Poly = vector<int>;
                                                                auto L = q[i << 1], R = q[i << 1 | 1];</pre>
                                                               q[i << 1 | 0] = MulT(q[i], R, L.size());
q[i << 1 | 1] = MulT(q[i], L, R.size());</pre>
auto Mul(auto a, auto b, size_t sz) {
 a.resize(sz), b.resize(sz);
 ntt(a.data(), sz); ntt(b.data(), sz);
                                                               for (int i = 0; i < n; i++) ans[i] = q[i + n][0];</pre>
 fi(0, sz) a[i] = modmul(a[i], b[i]);
 return ntt(a.data(), sz, true), a;
                                                               return ans.resize(x.size()), ans;
Poly Newton(const Poly &v, auto &&init, auto &&iter) {
                                                              pair<Poly, Poly> DivMod(const Poly &A, const Poly &B) {
                                                               assert(!B.empty() && B.back() != 0);
 Poly Q = { init(v[0]) };
                                                               if (A.size() < B.size()) return {{}}, A};</pre>
 for (int sz = 2; Q.size() < v.size(); sz *= 2) {</pre>
  Poly A{begin(v), begin(v) + min(sz, int(v.size()))};
                                                               const int sz = A.size() - B.size() + 1;
                                                               Poly X = B; reverse(all(X)); X.resize(sz);
Poly Y = A; reverse(all(Y)); Y.resize(sz);
  A.resize(sz \star 2), Q.resize(sz \star 2);
  iter(Q, A, sz * 2); Q.resize(sz);
                                                               Poly Q = Mul(Inv(X), Y);
                                                               Q.resize(sz); reverse(all(Q)); X = Mul(Q, B); Y = A;
 return Q.resize(v.size()), Q;
                                                               fi(0, Y.size()) Y[i] = modsub(Y[i], X[i]);
Poly Inv(const Poly &v) { // v[0] != 0
                                                               while (Y.size() && Y.back() == 0) Y.pop_back();
                                                               while (Q.size() && Q.back() == 0) Q.pop_back();
 return Newton(v, modinv,
  [](Poly &X, Poly &A, int sz) {
                                                               return {Q, Y};
   ntt(X.data(), sz), ntt(A.data(), sz);
for (int i = 0; i < sz; i++)</pre>
                                                              } // empty means zero polynomial
                                                              int LinearRecursionKth(Poly a, Poly c, int64_t k) {
    X[i] = modmul(X[i], modsub(2, modmul(X[i], A[i])));
                                                               const auto d = a.size(); assert(c.size() == d + 1);
                                                               const int sz = bit_ceil(2 * d + 1), o = sz / 2;
   ntt(X.data(), sz, true); });
                                                               Poly q = c; for (int &x: q) x = modsub(0, x); q[0]=1;
                                                               Poly p = Mul(a, q); p.resize(sz); q.resize(sz);
Poly Dx(Poly A) {
                                                               for (int r; r = (k & 1), k; k >>= 1) {
fi(1, A.size()) A[i - 1] = modmul(i, A[i]);
                                                                fill(d + all(p), 0); fill(d + 1 + all(q), 0);
 return A.empty() ? A : (A.pop_back(), A);
                                                                ntt(p.data(), sz); ntt(q.data(), sz);
                                                                for(int i = 0; i < sz; i++)
Poly Sx(Poly A) {
 A.insert(A.begin(), 0);
                                                                 p[i] = modmul(p[i], q[(i + o) & (sz - 1)]);
                                                                for(int i = 0, j = 0; j < sz; i++, j++)
q[i] = q[j] = modmul(q[i], q[j]);</pre>
 fi(1, A.size()) A[i] = modmul(modinv(i), A[i]);
 return A;
                                                                ntt(p.data(), sz, true); ntt(q.data(), sz, true);
Poly Ln(const Poly &A) { // coef[0] == 1; res[0] == 0
                                                                for (int i = 0; i < d; i++) p[i] = p[i << 1 | r];</pre>
                                                                for (int i = 0; i <= d; i++) q[i] = q[i << 1];
 auto B = Sx(Mul(Dx(A), Inv(A), bit_ceil(A.size()*2)));
                                                               } // Bostan-Mori
 return B.resize(A.size()), B;
                                                               return modmul(p[0], modinv(q[0]));
Poly Exp(const Poly &v) { // coef[0] == 0; res[0] == 1
                                                             return Newton(v, [](int x) { return 1; },
                                                                     Partition Number 19668451
  [](Poly &X, Poly &A, int sz) {
   auto Y = X; Y.resize(sz / 2); Y = Ln(Y);
                                                              ans[0] = tmp[0] = 1;
   fi(0, Y.size()) Y[i] = modsub(A[i], Y[i]);
                                                              for (int i = 1; i * i <= n; i++) {</pre>
   Y[0] = modadd(Y[0], 1); X = Mul(X, Y, sz); \});
                                                               for (int rep = 0; rep < 2; rep++)</pre>
                                                                for (int j = i; j <= n - i * i; j++)
modadd(tmp[j], tmp[j-i]);</pre>
Poly Pow(Poly a, lld M) { // period mod*(mod-1)
 assert(!a.empty() && a[0] != 0);
                                                               for (int j = i * i; j <= n; j++)</pre>
 const int N = int(a.size()); // mod x^N
                                                                modadd(ans[j], tmp[j - i * i]);
 const auto imul = [&a](int s) {
  for (int &x: a) x = modmul(x, s); }; int c = a[0];
                                                              5.18 Pi Count (+Linear Sieve) [8a4382]
 imul(modinv(c)); a = Ln(a); imul(M % mod);
                                                              static constexpr int N = 1000000 + 5;
 a = Exp(a); imul(modpow(c, M % (mod - 1)));
 return a;
                                                              lld pi[N]; vector<int> primes; bool sieved[N];
```

lld cube\_root(lld x) {

lld s = cbrt(x - 0.1L);

while (s \* s \* s <= x) ++s;

```
return s - 1;
lld square_root(lld x) {
lld s = sqrt(x - 0.1L);
while (s * s <= x) ++s;
return s - 1;
void init() {
primes.reserve(N);
for (int i = 2; i < N; i++) {</pre>
 if (!sieved[i]) primes.push_back(i);
 pi[i] = !sieved[i] + pi[i - 1];
  for (int p : primes) {
  if (i * p >= N) break;
   sieved[p * i] = true;
   if (i % p == 0) break;
 }
primes.insert(primes.begin(), 1);
lld phi(lld m, lld n) {
static constexpr int MM = 80000, NN = 500;
static lld val[MM][NN];
if (m<MM && n<NN && val[m][n]) return val[m][n] - 1;</pre>
 if (n == 0) return m;
 if (primes[n] >= m) return 1;
lld ret = phi(m, n - 1) - phi(m / primes[n], n - 1);
if (m < MM && n < NN) val[m][n] = ret + 1;</pre>
return ret:
lld pi_count(lld);
lld P2(lld m, lld n) {
lld sm = square_root(m), ret = 0;
for (lld i = n + 1; primes[i] <= sm; i++)</pre>
 ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
return ret;
lld pi_count(lld m) {
 if (m < N) return pi[m];</pre>
lld n = pi_count(cube_root(m));
return phi(m, n) + n - 1 - P2(m, n);
```

#### 5.19 Miller Rabin [fbd812]

```
bool isprime(llu x) {
auto witn = [&](llu a, int t) {
 for (llu a2; t--; a = a2) {
  a2 = mmul(a, a, x);
   if (a2 == 1 && a != 1 && a != x - 1) return true;
 }
 return a != 1;
if (x <= 2 || ~x & 1) return x == 2;
int t = countr_zero(x-1); llu odd = (x-1) >> t;
for (llu m:
 {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
 if (m % x != 0 && witn(mpow(m % x, odd, x), t))
  return false:
return true:
} // test @ luogu 143 & yosupo judge, ~1700ms for Q=1e5
// if use montgomery, ~250ms for Q=1e5
```

# 5.20 Pollard Rho [57ad88]

```
// does not work when n is prime or n == 1
  return any non-trivial factor
llu pollard_rho(llu n) {
 static mt19937_64 rnd(120821011);
 if (!(n & 1)) return 2;
 llu y = 2, z = y, c = rnd() % n, p = 1, i = 0, t;
 auto f = [&](llu x) {
  return madd(mmul(x, x, n), c, n); };
 do {
 p = mmul(msub(z = f(f(z)), y = f(y), n), p, n);
  if (++i &= 63) if (i == (i & -i)) t = gcd(p, n);
 } while (t == 1);
 return t == n ? pollard_rho(n) : t;
} // test @ yosupo judge, ~270ms for Q=100 // if use montgomery, ~70ms for Q=100
```

# Berlekamp Massey [a94d00]

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
```

```
vector<T> d(output.size() + 1), me, he;
for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
 for (size_t j = 0; j < me.size(); ++j)</pre>
  d[i] += output[i - j - 2] * me[j];
 if ((d[i] -= output[i - 1]) == 0) continue;
 if (me.empty()) {
  me.resize(f = i);
  continue;
 vector<T> o(i - f - 1);
 T k = -d[i] / d[f]; o.push_back(-k);
 for (T x : he) o.push_back(x * k);
 if (o.size() < me.size()) o.resize(me.size());</pre>
 for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
 if (i-f+he.size() >= me.size()) he = me, f = i;
 me = o;
return me:
```

### 5.22 Gauss Elimination [9dea40]

```
void gauss(vector<vector<llf>> A, vector<llf>> b) {
 const int n = A.size(), m = A[0].size();
 for (int i = 0; i < m; ++i) {</pre>
  int p = -1;
  for (int j = i; j < n; ++j)</pre>
   if (abs(A[j][i]) > eps)
    if (p == -1 || abs(A[j][i]) > abs(A[p][i]))
     p = j;
  if (p == -1) continue;
  swap(A[p], A[i]);
  swap(b[p], b[i]);
  for (int j = 0; j < n; ++j) if (j != i) {</pre>
   llf z = A[j][i] / A[i][i];
   for (int k = 0; k < m; ++k)
    A[j][k] = z * A[i][k];
   b[j] -= z * b[i];
}
```

# Charateristic Polynomial [ff2159]

```
#define rep(x, y, z) for (int x=y; x<z; x++)
using VI = vector<int>; using VVI = vector<VI>;
void Hessenberg(VVI &H, int N) {
   for (int i = 0; i < N - 2; ++i) {</pre>
  for (int j = i + 1; j < N; ++j) if (H[j][i]) {
   rep(k, i, N) swap(H[i+1][k], H[j][k]);
   rep(k, 0, N) swap(H[k][i+1], H[k][j]);
   break;
  if (!H[i + 1][i]) continue;
  for (int j = i + 2; j < N; ++j) {
   int co = mul(modinv(H[i + 1][i]), H[j][i]);
   rep(k, i, N) subeq(H[j][k], mul(H[i+1][k], co));
   rep(k, 0, N) addeq(H[k][i+1], mul(H[k][j], co));
 }
VI CharacteristicPoly(VVI &A) {
 int N = (int)A.size(); Hessenberg(A, N);
 VVI P(N + 1, VI(N + 1)); P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {</pre>
  rep(j, 0, i+1) P[i][j] = j ? P[i-1][j-1] : 0;
  for (int j = i - 1, val = 1; j >= 0; --j) {
   int co = mul(val, A[j][i - 1]);
   rep(k, 0, j+1) subeq(P[i][k], mul(P[j][k], co));
   if (j) val = mul(val, A[j][j - 1]);
 if (N \& 1) for (int \&x: P[N]) x = sub(0, x);
 return P[N]; // test: 2021 PTZ Korea K
```

# **5.24** Simplex [c9c93b]

```
namespace simplex {
// maximize c^Tx under Ax \le B and x \ge 0
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
using VD = vector<llf>;
using VVD = vector<vector<llf>>;
const llf eps = 1e-9, inf = 1e+9;
```

```
int n, m; VVD d; vector<int> p, q;
void pivot(int r, int s) {
llf inv = 1.0 / d[r][s];
 for (int i = 0; i < m + 2; ++i)</pre>
  for (int j = 0; j < n + 2; ++j)
   if (i != r && j != s)
    d[i][j] -= d[r][j] * d[i][s] * inv;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;</pre>
 for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
 d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
 int x = m + z;
 while (true) {
  int s = -1;
  for (int i = 0; i <= n; ++i) {</pre>
   if (!z && q[i] == -1) continue;
   if (s == -1 || d[x][i] < d[x][s]) s = i;
  if (s == -1 || d[x][s] > -eps) return true;
  int r = -1;
  for (int i = 0; i < m; ++i) {
   if (d[i][s] < eps) continue;
if (r == -1 ||</pre>
    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
  pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
m = (int)b.size(), n = (int)c.size();
 d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)
 for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)
  p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
 for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];</pre>
 q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
                                                                       }
 for (int i = 1; i < m; ++i)</pre>
  if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
 if (d[r][n + 1] < -eps) {</pre>
  pivot(r, n);
  if (!phase(1) || d[m + 1][n + 1] < -eps)
   return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
   int s = min_element(d[i].begin(), d[i].end() - 1)
        - d[i].begin();
   pivot(i, s);
  }
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)</pre>
  if (p[i] < n) x[p[i]] = d[i][n + 1];</pre>
 return x;
}}
5.25 Simplex Construction
Standard form: maximize \sum_{1 \le i \le n} c_i x_i such that for all 1 \le j \le m, 6.3 2D Farthest Pair [8b5844]
\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j and x_i \geq 0 for all 1 \leq i \leq n.
  \bar{1}. In case of minimization, let c_i'=-c_i
  2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
  3. \sum_{1 \le i \le n}^{-} A_{ji} x_i = b_j
        • \sum_{1 \le i \le n} A_{ji} x_i \le b_j
        • \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j
     If x_i has no lower bound, replace x_i with x_i
5.26 Adaptive Simpson [09669e]
llf simp(llf l, llf r) {
 llf m = (l + r) / 2;
 return (f(l) + f(r) + 4.0 * f(m)) * (r - l) / 6.0;
ilf F(llf L, llf R, llf v, llf eps) {
    llf M = (L + R) / 2, vl = simp(L, M), vr = simp(M, R);
    if (abs(vl + vr - v) <= 15 * eps)</pre>
  return vl + vr + (vl + vr - v) / 15.0;
```

**return** F(L, M, vl, eps / 2.0) +

F(M, R, vr, eps / 2.0);

} // call F(l, r, simp(l, r), 1e-6)

```
Geometry
```

# Basic Geometry [e4a147]

```
#define IM imag
#define RE real
using lld = int64_t;
using llf = long double;
using PT = std::complex<lld>;
using PTF = std::complex<llf>;
using P = PT;
llf abs(P p) { return sqrtl(norm(p)); }
PTF toPTF(PT p) { return PTF{RE(p), IM(p)}; }
int sgn(lld x) { return (x > 0) - (x < 0); }</pre>
lld dot(P a, P b) { return RE(conj(a) * b); }
lld cross(P a, P b) { return IM(conj(a) * b); }
int ori(P a, P b, P c) {
return sgn(cross(b - a, c - a));
int quad(P p) {
 return (IM(p) == 0) // use sgn for PTF
  ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
int argCmp(P a, P b) {
 // returns 0/+-1, starts from theta = -PI
 int qa = quad(a), qb = quad(b);
 if (qa != qb) return sgn(qa - qb);
 return sgn(cross(b, a));
P rot90(P p) { return P{-IM(p), RE(p)}; }
template <typename V> llf area(const V & pt) {
 lld ret = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 return ret / 2.0;
template <typename V> PTF center(const V & pt) {
 P ret = 0; lld A = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++) {</pre>
  lld cur = cross(pt[i] - pt[0], pt[i+1] - pt[0]);
  ret += (pt[i] + pt[i + 1] + pt[0]) * cur; A += cur;
 return toPTF(ret) / llf(A * 3);
PTF project(PTF p, PTF q) { // p onto q
 return dot(p, q) * q / dot(q, q); // dot<llf>
```

#### 6.2 2D Convex Hull [ecba37]

```
// from NaCl, counterclockwise, be careful of n<=2
vector<P> convex_hull(vector<P> v) {
 sort(all(v)); // by X then Y
 if (v[0] == v.back()) return {v[0]};
 int t = 0, s = 1; vector<P> h(v.size() + 1);
 for (int _ = 2; _--; s = t--, reverse(all(v)))
  for (P p : v) {
   while (t>s && ori(p, h[t-1], h[t-2]) >= 0) t--;
  h[t++] = p;
 return h.resize(t), h;
```

```
// p is CCW convex hull w/o colinear points
int n = (int)p.size(), pos = 1; lld ans = 0;
for (int i = 0; i < n; i++) {</pre>
 P = p[(i + 1) \% n] - p[i];
 while (cross(e, p[(pos + 1) % n] - p[i]) >
      cross(e, p[pos] - p[i]))
  pos = (pos + 1) % n;
 for (int j: {i, (i + 1) % n})
ans = max(ans, norm(p[pos] - p[j]));
} // tested @ AOJ CGL_4_B
```

# 6.4 MinMax Enclosing Rect [e4470c]

```
// from 8BQube, plz ensure p is strict convex hull
const llf INF = 1e18, qi = acos(-1) / 2 * 3;
pair<llf, llf> solve(const vector<P> &p) {
llf mx = 0, mn = INF; int n = (int)p.size();
 for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {
#define Z(v) (p[(v) % n] - p[i])
  P = Z(i + 1);
  while (cross(e, Z(u + 1)) > cross(e, Z(u))) ++u;
```

```
while (dot(e, Z(r + 1)) > dot(e, Z(r))) ++r;
if (!i) l = r + 1;
while (dot(e, Z(l + 1)) < dot(e, Z(l))) ++l;
P D = p[r % n] - p[l % n];
llf H = cross(e, Z(u)) / llf(norm(e));
mn = min(mn, dot(e, D) * H);
llf B = sqrt(norm(D)) * sqrt(norm(Z(u)));
llf deg = (qi - acos(dot(D, Z(u)) / B)) / 2;
mx = max(mx, B * sin(deg) * sin(deg));
}
return {mn, mx};
} // test @ UVA 819</pre>
```

# 6.5 Minkowski Sum [602806]

```
// A, B are strict convex hull rotate to min by (X, Y)
vector<P> Minkowski(vector<P> A, vector<P> B) {
    const int N = (int)A.size(), M = (int)B.size();
    vector<P> sa(N), sb(M), C(N + M + 1);
    for (int i = 0; i < N; i++) sa[i] = A[(i+1)%N]-A[i];
    for (int i = 0; i < M; i++) sb[i] = B[(i+1)%M]-B[i];
    C[0] = A[0] + B[0];
    for (int i = 0, j = 0; i < N || j < M; ) {
        P e = (j>=M || (i<N && cross(sa[i], sb[j])>=0))
        ? sa[i++] : sb[j++];
    C[i + j] = e;
    }
    partial_sum(all(C), C.begin()); C.pop_back();
    return convex_hull(C); // just to remove colinear
}
```

### 6.6 Segment Intersection [60d016]

```
struct Seg { // closed segment
P st, dir; // represent st + t*dir for 0<=t<=1
Seg(P s, P e) : st(s), dir(e - s) {}
static bool valid(lld p, lld q) {
  // is there t s.t. 0 <= t <= 1 && qt == p ?
 if (q < 0) q = -q, p = -p;
  return 0 <= p && p <= q;
vector<P> ends() const { return { st, st + dir }; }
template <typename T> bool isInter(T A, P p) {
if (A.dir == P(0)) return p == A.st; // BE CAREFUL
return cross(p - A.st, A.dir) == 0 &&
 T::valid(dot(p - A.st, A.dir), norm(A.dir));
template <typename U, typename V>
bool isInter(U A, V B) {
 if (cross(A.dir, B.dir) == 0) { // BE CAREFUL
 bool res = false;
 for (P p: A.ends()) res |= isInter(B, p);
  for (P p: B.ends()) res |= isInter(A, p);
  return res;
P D = B.st - A.st; lld C = cross(A.dir, B.dir);
return U::valid(cross(D, B.dir), C) &&
  V::valid(cross(D, A.dir), C);
```

#### 6.7 Half Plane Intersection [45e909]

```
struct Line {
P st, ed, dir;
Line (P s, P e) : st(s), ed(e), dir(e - s) {}
}; using LN = const Line &;
PTF intersect(LN A, LN B) {
llf t = cross(B.st - A.st, B.dir) /
 llf(cross(A.dir, B.dir));
return toPTF(A.st) + toPTF(A.dir) * t; // C^3 / C^2
bool cov(LN l, LN A, LN B) {
i128 u = cross(B.st-A.st, B.dir);
i128 v = cross(A.dir, B.dir);
// ori(l.st, l.ed, A.st + A.dir*(u/v)) <= 0?
i128 x = RE(A.dir) * u + RE(A.st - l.st) * v;
i128 y = IM(A.dir) * u + IM(A.st - l.st) * v;
return sgn(x*IM(l.dir) - y*RE(l.dir)) * sgn(v) >= 0;
} // x, y are C^3, also sgn<i128> is needed
bool operator<(LN a, LN b) {</pre>
if (int c = argCmp(a.dir, b.dir)) return c == -1;
return ori(a.st, a.ed, b.st) < 0;</pre>
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
```

```
/ the half plane is the LHS when going from st to ed
llf HPI(vector<Line> &q) {
 sort(q.begin(), q.end());
 int n = (int)q.size(), l = 0, r = -1;
for (int i = 0; i < n; i++) {</pre>
  if (i && !argCmp(q[i].dir, q[i-1].dir)) continue;
  while (l < r && cov(q[i], q[r-1], q[r])) --r;
while (l < r && cov(q[i], q[l], q[l+1])) ++l;</pre>
  q[++r] = q[i];
 while (l < r && cov(q[l], q[r-1], q[r])) --r;</pre>
 while (l < r && cov(q[r], q[l], q[l+1])) ++l;</pre>
 n = r - l + 1; // q[l .. r] are the lines
 if (n <= 1 || !argCmp(q[l].dir, q[r].dir)) return 0;</pre>
 vector<PTF> pt(n);
 for (int i = 0; i < n; i++)</pre>
  pt[i] = intersect(q[i+l], q[(i+1)%n+l]);
 return area(pt);
} // test @ 2020 Nordic NCPC : BigBrother
6.8 SegmentDist (Sausage) [9d8603]
// be careful of abs<complex<int>> (replace _abs below)
```

```
// be careful of abs<complex<int>> (replace _abs below
llf PointSegDist(P A, Seg B) {
   if (B.dir == P(0)) return _abs(A - B.st);
   if (sgn(dot(A - B.st, B.dir)) *
      sgn(dot(A - B.ed, B.dir)) <= 0)
      return abs(cross(A - B.st, B.dir)) / _abs(B.dir);
   return min(_abs(A - B.st), _abs(A - B.ed));
}
llf SegSegDist(const Seg &s1, const Seg &s2) {
   if (isInter(s1, s2)) return 0;
   return min({
      PointSegDist(s1.st, s2),
      PointSegDist(s2.st, s1),
      PointSegDist(s2.ed, s1) });
} // test @ QOJ2444 / PTZ19 Summer.D3</pre>
```

# 6.9 Rotating Sweep Line [8aff27]

```
struct Event {
 Pd; int u, v;
 bool operator<(const Event &b) const {</pre>
  return sgn(cross(d, b.d)) > 0; }
P makePositive(P z) { return cmpxy(z, 0) ? -z : z; }
void rotatingSweepLine(const vector<P> &p) {
 const int n = int(p.size());
 vector<Event> e; e.reserve(n * (n - 1) / 2);
 for (int i = 0; i < n; i++)</pre>
  for (int j = i + 1; j < n; j++)
    e.emplace_back(makePositive(p[i] - p[j]), i, j);
 sort(all(e));
 vector<int> ord(n), pos(n);
 iota(all(ord), 0);
sort(all(ord), [&p](int i, int j) {
   return cmpxy(p[i], p[j]); });
 for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
 const auto makeReverse = [](auto &v) {
  sort(all(v)); v.erase(unique(all(v)), v.end());
  vector<pair<int,int>> segs;
  for (size_t i = 0, j = 0; i < v.size(); i = j) {
   for (; j < v.size() && v[j] - v[i] <= j - i; j++);
   segs.emplace_back(v[i], v[j - 1] + 1 + 1);</pre>
  }
  return segs;
 for (size_t i = 0, j = 0; i < e.size(); i = j) {</pre>
  /* do here */
  vector<size_t> tmp;
  for (; j < e.size() && !(e[i] < e[j]); j++)</pre>
   tmp.push_back(min(pos[e[j].u], pos[e[j].v]));
  for (auto [l, r] : makeReverse(tmp)) {
   reverse(ord.begin() + l, ord.begin() + r);
    for (int t = l; t < r; t++) pos[ord[t]] = t;</pre>
 }
}
```

#### 6.10 Polygon Cut [e9bdd1]

```
using P = PTF;
vector<P> cut(const vector<P>& poly, P s, P e) {
  vector<P> res;
```

```
National Taiwan University - ckiseki
 for (size_t i = 0; i < poly.size(); i++) {</pre>
  P cur = poly[i], prv = i ? poly[i-1] : poly.back();
                                                                assert (succ);
  bool side = ori(s, e, cur) < 0;</pre>
                                                               } else {
  if (side != (ori(s, e, prv) < 0))
                                                                int qid = i - N;
   res.push_back(intersect({s, e}, {cur, prv}));
  if (side)
                                                                 ans[qid] = 1;
   res.push_back(cur);
                                                                 continue;
 return res;
}
6.11 Point In Simple Polygon [037c52]
bool PIP(const vector<P> &p, P z, bool strict = true) {
                                                                 continue;
 int cnt = 0, n = (int)p.size();
 for (int i = 0; i < n; i++) {</pre>
  P A = p[i], B = p[(i + 1) \% n];
                                                                 ) {
  if (isInter(Seg(A, B), z)) return !strict;
  auto zy = IM(z), Ay = IM(A), By = IM(B);
                                                                 continue:
 cnt ^= ((zy<Ay) - (zy<By)) * ori(z, A, B) > 0;
 return cnt:
}
6.12 Point In Hull (Fast) [060ba1]
bool PIH(const vector<P> &h, P z, bool strict = true) {
                                                              return ans;
 int n = (int)h.size(), a = 1, b = n - 1, r = !strict;
 if (n < 3) return r && isInter(Seg(h[0], h[n-1]), z);</pre>
 if (ori(h[0],h[a],h[b]) > 0) swap(a, b);
 if (ori(h[0],h[a],z) >= r || ori(h[0],h[b],z) <= -r)</pre>
  return false;
 while (abs(a - b) > 1) {
  int c = (a + b) / 2;
  (ori(h[0], h[c], z) > 0 ? b : a) = c;
 return ori(h[a], h[b], z) < r;
                                                               else l = m:
6.13 Point In Polygon (Fast) [71725b]
vector<int> PIPfast(vector<P> p, vector<P> q) {
 const int N = int(p.size()), Q = int(q.size());
 vector<pair<P, int>> evt;
 vector<Seg> edge;
 for (int i = 0; i < N; i++) {</pre>
  int a = i, b = (i + 1) % N;
 P A = p[a], B = p[b];
                                                            llf FMOD(llf x) {
  assert (A < B || B < A); // std::operator<</pre>
 if (B < A) swap(A, B);
 evt.emplace_back(A, i);
evt.emplace_back(B, ~i);
                                                             return x;
 edge.emplace_back(A, B);
 for (int i = 0; i < Q; i++)</pre>
 evt.emplace_back(q[i], i + N);
 sort(all(evt));
```

```
auto vtx = p; sort(all(vtx));
auto eval = [](const Seg &a, lld x) -> llf {
 if (RE(a.dir) == 0) {
  assert (x == RE(a.st));
  return IM(a.st) + llf(IM(a.dir)) / 2;
 llf t = (x - RE(a.st)) / llf(RE(a.dir));
 return IM(a.st) + IM(a.dir) * t;
lld cur_x = 0;
auto cmp = [&](const Seg &a, const Seg &b) -> bool {
 if (int s = sgn(eval(a, cur_x) - eval(b, cur_x)))
  return s == -1;
 int s = sgn(cross(b.dir, a.dir));
 if (cur_x != RE(a.st) && cur_x != RE(b.st)) s *= -1;
 return s == -1;
};
namespace pbds = __gnu_pbds;
using Tree = pbds::tree<Seg, int, decltype(cmp),</pre>
   pbds::rb_tree_tag,
   pbds::tree_order_statistics_node_update>;
Tree st(cmp);
vector<int> ans(Q);
for (auto [ep, i] : evt) {
 cur_x = RE(ep);
 if (i < 0) { // remove
  st.erase(edge[~i]);
 } else if (i < N) { // insert</pre>
```

```
17
   auto [it, succ] = st.insert({edge[i], i});
   if (binary_search(all(vtx), ep)) { // on vertex
   Seg H(ep, ep); // ??
   auto it = st.lower_bound(H);
   if (it != st.end() && isInter(it->first, ep)) {
    ans[qid] = 1; // on edge
   if (it != st.begin() && isInter(prev(it)->first, ep)
    ans[qid] = 1; // on edge
   auto rk = st.order_of_key(H);
   if (rk % 2 == 0) ans[qid] = 0; // outside
   else ans[qid] = 2; // inside
} // test @ AOJ CGL_3_C
6.14 Tangent of Points To Hull [6d7cd7]
 const auto gao = [&, N = int(v.size())](int s) {
  const auto lt = [&](int x, int y) {
  return ori(p, v[x % N], v[y % N]) == s; };
  int l = 0, r = N; bool up = lt(0, 1);
  while (r - l > 1) {
   int m = (l + r) / 2;
```

```
pair<int, int> get_tangent(const vector<P> &v, P p) {
   if (lt(m, 0) ? up : !lt(m, m+1)) r = m;
  return (lt(l, r) ? r : l) % N;
 }; // test @ codeforces.com/gym/101201/problem/E
 return {gao(-1), gao(1)}; // (a,b):ori(p,v[a],v[b])<0</pre>
} // plz ensure that point strictly out of hull
```

# 6.15 Circle Class & Intersection [5111af]

```
if (x < -PI) x += PI * 2;
 if (x > PI) x -= PI * 2;
struct Cir { PTF o; llf r; };
// be carefule when tangent
vector<llf> intersectAngle(Cir a, Cir b) {
PTF dir = b.o - a.o; llf d2 = norm(dir);
 if (norm(a.r - b.r) >= d2) { // <math>norm(x) := |x|^2}
 if (a.r < b.r) return {-PI, PI}; // a in b</pre>
  else return {}; // b in a
 } else if (norm(a.r + b.r) <= d2) return {};</pre>
 llf dis = abs(dir), theta = arg(dir);
 llf phi = acos((a.r * a.r + d2 - b.r * b.r) /
   (2 * a.r * dis)); // is acos_safe needed ?
 llf L = FMOD(theta - phi), R = FMOD(theta + phi);
 return { L, R };
vector<PTF> intersectPoint(Cir a, Cir b) {
 llf d = abs(a.o - b.o);
 if (d > b.r+a.r || d < abs(b.r-a.r)) return {};</pre>
 llf dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
 PTF dir = (a.o - b.o) / d;
PTF u = dir * d1 + b.o;
 PTF v = rot90(dir) * sqrt(max(0.0L, b.r*b.r-d1*d1));
 return \{u + v, u - v\};
} // test @ AOJ CGL probs
```

#### 6.16 Circle Common Tangent [5ff02c]

```
// be careful of tangent / exact same circle
// sign1 = 1 for outer tang, -1 for inter tang
vector<Line> common_tan(const Cir &a, const Cir &b, int
     sign1) {
 if (norm(a.o - b.o) < eps) return {};</pre>
llf d = abs(a.o - b.o), c = (a.r - sign1 * b.r) / d;
 PTF v = (b.o - a.o) / d;
if (c * c > 1) return {};
```

```
if (abs(c * c - 1) < eps) {
   PTF p = a.o + c * v * a.r;
   return {Line(p, p + rot90(b.o - a.o))};
}
vector<Line> ret; llf h = sqrt(max(0.0L, 1-c*c));
for (int sign2 : {1, -1}) {
   PTF n = c * v + sign2 * h * rot90(v);
   PTF p1 = a.o + n * a.r;
   PTF p2 = b.o + n * (b.r * sign1);
   ret.emplace_back(p1, p2);
}
return ret;
}
```

#### 6.17 Line-Circle Intersection [12b42a]

```
vector<PTF> LineCircleInter(PTF p1, PTF p2, PTF o, llf
    r) {
    PTF ft = p1 + project(o-p1, p2-p1), vec = p2-p1;
    llf dis = abs(o - ft);
    if (abs(dis - r) < eps) return {ft};
    if (dis > r) return {};
    vec = vec * sqrt(r * r - dis * dis) / abs(vec);
    return {ft + vec, ft - vec}; // sqrt_safe?
}
```

### 6.18 Poly-Circle Intersection [7f140a]

```
// Divides into multiple triangle, and sum up
// from 8BQube, test by HDU2892 & AOJ CGL_7_H
llf _area(PTF pa, PTF pb, llf r) {
  if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
 if (abs(pb) < eps) return 0;</pre>
 llf S, h, theta;
 llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
 llf cB = dot(pb, pb-pa) / a / c, B = acos_safe(cB);
llf cC = dot(pa, pb) / a / b, C = acos_safe(cC);
 if (a > r) {
  S = (C / 2) * r * r; h = a * b * sin(C) / c;
  if (h < r && B < PI / 2)
   S = (acos\_safe(h/r)*r*r - h*sqrt\_safe(r*r-h*h));
 } else if (b > r) {
  theta = PI - B - asin_safe(sin(B) / r * a);
  S = 0.5 * a*r*sin(theta) + (C-theta)/2 * r * r;
 } else
  S = 0.5 * sin(C) * a * b;
 return S;
llf area_poly_circle(const vector<PTF> &v, PTF 0, llf r
    ) {
 llf S = 0;
 for (size_t i = 0, N = v.size(); i < N; ++i)</pre>
  S += _area(v[i] - 0, v[(i + 1) % N] - 0, r) *
     ori(0, v[i], v[(i + 1) % N]);
 return abs(S);
}
```

#### 6.19 Minimum Covering Circle [faa85a]

```
Cir getCircum(P a, P b, P c){ // P = complex<llf>
 P z1 = a - b, z2 = a - c; llf D = cross(z1, z2) * 2;
 llf c1 = dot(a + b, z1), c2 = dot(a + c, z2);
 P \circ = rot90(c2 * z1 - c1 * z2) / D;
return { o, abs(o - a) };
Cir minCircleCover(vector<P> pts) {
assert (!pts.empty());
 ranges::shuffle(pts, mt19937(114514));
 Cir c = { 0, 0 };
for(size_t i = 0; i < pts.size(); i++) {</pre>
 if (abs(pts[i] - c.o) <= c.r) continue;</pre>
  c = { pts[i], 0 };
for (size_t j = 0; j < i; j++) {</pre>
   if (abs(pts[j] - c.o) <= c.r) continue;</pre>
   c.o = (pts[i] + pts[j]) / llf(2);
   c.r = abs(pts[i] - c.o);
   for (size_t k = 0; k < j; k++) {</pre>
    if (abs(pts[k] - c.o) <= c.r) continue;</pre>
    c = getCircum(pts[i], pts[j], pts[k]);
   }
  }
 return c;
} // test @ TIOJ 1093 & luogu P1742
```

### 6.20 Circle Union [1a5265]

```
#define eb emplace_back
struct Teve { // test@SPOJ N=1000, 0.3~0.5s
 PTF p; llf a; int add; // point, ang, add
 Teve(PTF x, llf y, int z) : p(x), a(y), add(z) {}
 bool operator<(Teve &b) const { return a < b.a; }</pre>
// strict: x = 0, otherwise x = -1
bool disjunct(Cir &a, Cir &b, int x)
{ return sgn(abs(a.o - b.o) - a.r - b.r) > x; }
bool contain(Cir &a, Cir &b, int x)
{ return sgn(a.r - b.r - abs(a.o - b.o)) > x; }
vector<llf> CircleUnion(vector<Cir> &c) {
 // area[i] : area covered by at least i circles
 int N = (int)c.size(); vector<llf> area(N + 1);
 vector<vector<int>> overlap(N, vector<int>(N));
 auto g = overlap; // use simple 2darray to speedup
 for (int i = 0; i < N; ++i)</pre>
  for (int j = 0; j < N; ++j) {
   /* c[j] is non-strictly in c[i]. */
   overlap[i][j] = i != j &&
    (sgn(c[i].r - c[j].r) > 0 | |
      (sgn(c[i].r - c[j].r) == 0 && i < j)) &&
    contain(c[i], c[j], -1);
 for (int i = 0; i < N; ++i)</pre>
  for (int j = 0; j < N; ++j)
g[i][j] = i != j && !(overlap[i][j] ||</pre>
     overlap[j][i] || disjunct(c[i], c[j], -1));
 for (int i = 0; i < N; ++i) {</pre>
  vector<Teve> eve; int cnt = 1;
  for (int j = 0; j < N; ++j) cnt += overlap[j][i];</pre>
  // if (cnt > 1) continue; (if only need area[1])
  for (int j = 0; j < N; ++j) if (g[i][j]) {</pre>
   auto IP = intersectPoint(c[i], c[j]);
   PTF aa = IP[1], bb = IP[0];
   llf A = arg(aa - c[i].o), B = arg(bb - c[i].o);
   eve.eb(bb, B, 1); eve.eb(aa, A, -1);
   if (B > A) ++cnt;
  if (eve.empty()) area[cnt] += PI*c[i].r*c[i].r;
   sort(eve.begin(), eve.end());
   eve.eb(eve[0]); eve.back().a += PI * 2;
   for (size_t j = 0; j + 1 < eve.size(); j++) {</pre>
    cnt += eve[j].add;
    area[cnt] += cross(eve[j].p, eve[j+1].p) *.5;
    llf t = eve[j + 1].a - eve[j].a;
    area[cnt] += (t-sin(t)) * c[i].r * c[i].r *.5;
  }
 }
 return area;
```

#### 6.21 Polygon Union [2bff43]

```
llf rat(P a, P b) { return sgn(RE(b)) ? llf(RE(a))/RE(b
    ) : llf(IM(a))/IM(b); }
llf polyUnion(vector<vector<P>>& poly) {
 llf ret = 0; // area of poly[i] must be non-negative
 rep(i,0,sz(poly)) rep(v,0,sz(poly[i])) {
  P A = poly[i][v], B = poly[i][(v + 1) % sz(poly[i])];
  vector<pair<llf, int>> segs{{0, 0}, {1, 0}};
  rep(j,0,sz(poly)) if (i != j) {
   rep(u,0,sz(poly[j])) {
    P C = poly[j][u], D = poly[j][(u + 1) % sz(poly[j])
    if (int sc = ori(A, B, C), sd = ori(A, B, D); sc !=
     sd) {
     llf sa = cross(D-C, A-C), sb = cross(D-C, B-C);
     if (min(sc, sd) < 0)
      segs.emplace_back(sa / (sa - sb), sgn(sc - sd));
    } else if (!sc && !sd && j<i && sgn(dot(B-A,D-C))</pre>
    >0){
     segs.emplace_back(rat(C - A, B - A), 1);
     segs.emplace_back(rat(D - A, B - A), -1);
    }
  }
  sort(segs.begin(), segs.end());
  for (auto &s : segs) s.first = clamp<llf>(s.first, 0,
     1);
```

```
National Taiwan University - ckiseki
  llf sum = 0;
                                                                  area += abs(ver(p[a], p[b], p[c]))/2.0,
  int cnt = segs[0].second;
                                                              // vol += volume(P3(0, 0, 0), p[a], p[b], p[c])/6.0;
  rep(j,1,sz(segs)) {
                                                              6.24 3D Projection [68f350]
   if (!cnt) sum += segs[j].first - segs[j - 1].first;
   cnt += segs[j].second;
                                                              using P3F = valarray<llf>;
                                                              P3F toP3F(P3 p) { return {p.x, p.y, p.z}; }
 }
 ret += cross(A,B) * sum;
                                                              llf dot(P3F a, P3F b) {
 return ret / 2;
                                                              P3F housev(P3 A, P3 B, int s) {
6.22 3D Point [46b73b]
struct P3 {
 lld x, y, z;
                                                              P project(P3 p, P3 q) {
 P3 operator^(const P3 &b) const {
                                                               P3 o(0, 0, 1);
  return {y*b.z-b.y*z, z*b.x-b.z*x, x*b.y-b.x*y};
                                                               auto pf = toP3F(p);
 //Azimuthal angle (longitude) to x-axis. \in [-pi, pi]
 llf phi() const { return atan2(y, x); }
                                                               return P(np[0], np[1]);
 //Zenith angle (latitude) to the z-axis. \in [0, pi]
 llf theta() const { return atan2(sqrt(x*x+y*y),z); }
P3 ver(P3 a, P3 b, P3 c) { return (b - a) ^ (c - a); }
                                                                 • L_1: v_1 = p_1 + t_1 d_1, L_2: v_2 = p_2 + t_2 d_2
                                                                 \cdot \; \boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2
lld volume(P3 a, P3 b, P3 c, P3 d) {
                                                                 • \boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}, \boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}
return dot(ver(a, b, c), d - a);
P3 rotate_around(P3 p, llf angle, P3 axis) {
                                                              6.26 Delaunay [3a4ff1]
llf s = sin(angle), c = cos(angle);
 P3 u = normalize(axis);
 return u*dot(u, p)*(1-c) + p * c + cross(u, p)*s;
6.23 3D Convex Hull [01652a]
                                                              #define L(i) ((i)==0 ? 2 : (i)-1)
struct Face {
                                                              #define R(i) ((i)==2 ? 0 : (i)+1)
 int a, b, c;
 Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
auto preprocess(const vector<P3> &pt) {
                                                               i128 inf_det = 0, det = 0, inf_N, N;
                                                               F3 {
 auto G = pt.begin();
 auto a = find_if(all(pt), [&](P3 z) {
  return z != *G; }) - G;
 auto b = find_if(all(pt), [&](P3 z) {
  return ver(*G, pt[a], z) != P3(0, 0, 0); }) - G;
 auto c = find_if(all(pt), [&](P3 z) {
  return volume(*G, pt[a], pt[b], z) != 0; }) - G;
                                                                inf_det += inf_N * D; det += N * D;
 vector<size_t> id;
 for (size_t i = 0; i < pt.size(); i++)</pre>
 if (i != a && i != b && i != c) id.push_back(i);
 return tuple{a, b, c, id};
                                                              P v[maxn];
                                                              struct Tri;
                                                              struct E {
// return the faces with pt indexes
// all points coplanar case will WA
                                                               Tri *t; int side;
vector<Face> convex_hull_3D(const vector<P3> &pt) {
 const int n = int(pt.size());
 if (n <= 3) return {}; // be careful about edge case</pre>
                                                              struct Tri {
 vector<Face> now;
 vector<vector<int>> z(n, vector<int>(n));
 auto [a, b, c, ord] = preprocess(pt);
 now.emplace_back(a, b, c); now.emplace_back(c, b, a);
                                                               bool contains(int q) const {
 for (auto i : ord) {
                                                                 return false;
  vector<Face> next;
  for (const auto &f : now) {
                                                                return true;
   lld v = volume(pt[f.a], pt[f.b], pt[f.c], pt[i]);
   if (v <= 0) next.push_back(f);</pre>
```

z[f.a][f.b] = z[f.b][f.c] = z[f.c][f.a] = sgn(v);

//  $n^2$  delaunay: facets with negative z normal of // convexhull of  $(x, y, x^2 + y^2)$ , use a pseudo-point

// llf area = 0, vol = 0; // surface area / volume

const auto  $F = [\&](int x, int y) \{$ 

**if** (z[x][y] > 0 && z[y][x] <= 0)

F(f.a, f.b), F(f.b, f.c), F(f.c, f.a);

// (0, 0, inf) to avoid degenerate case

next.emplace\_back(x, y, i);

for (const auto &f : now)

// for (auto [a, b, c]: faces)

now = next;

// test @ SPOJ CH3D

return now;

}

}

```
return a[0]*b[0]+a[1]*b[1]+a[2]*b[2];
 const llf a = abs(A), b = abs(B);
return toP3F(A) / a + s * toP3F(B) / b;
 P3F u = housev(q, o, q.z > 0 ? 1 : -1);
 auto np = pf - 2 * u * dot(u, pf) / dot(u, u);
} // project p onto the plane q^Tx = 0
6.25 3D Skew Line Nearest Point
   m{\cdot} \ \ m{c}_1 = m{p}_1 + rac{(m{p}_2 - m{p}_1) \cdot m{n}_2}{m{d}_1 \cdot m{n}_2} m{d}_1, m{c}_2 = m{p}_2 + rac{(m{p}_1 - m{p}_2) \cdot m{n}_1}{m{d}_2 \cdot m{n}_1} m{d}_2
/* please ensure input points are unique */
/* A triangulation such that no points will strictly
inside circumcircle of any triangle. C should be big
enough s.t. the initial triangle contains all points */
#define F3 for (int i = 0; i < 3; i++)
bool is_inf(P z) { return RE(z) \leftarrow C \mid \mid RE(z) >= C; }
bool in_cc(const array<P,3> &p, P q) {
  if (is_inf(p[i]) && is_inf(q)) continue;
  else if (is_inf(p[i])) inf_N = 1, N = -norm(q);
  else if (is_inf(q)) inf_N = -1, N = norm(p[i]);
  else inf_N = 0, N = norm(p[i]) - norm(q);
  lld D = \overline{cross(p[R(i)] - q, p[L(i)] - q)};
 return inf_det != 0 ? inf_det > 0 : det > 0;
 E(Tri *t_=0, int side_=0) : t(t_), side(side_) {}
 array<int,3> p; array<Tri*,3> ch; array<E,3> e;
 Tri(int a=0, int b=0, int c=0) : p{a, b, c}, ch{} {}
 bool has_chd() const { return ch[0] != nullptr; }
  F3 if (ori(v[p[i]], v[p[R(i)]], v[q]) < 0)
 bool check(int q) const {
  return in_cc({v[p[0]], v[p[1]], v[p[2]]}, v[q]); }
} pool[maxn * 10], *it, *root;
void link(const E &a, const E &b) {
 if (a.t) a.t->e[a.side] = b;
 if (b.t) b.t->e[b.side] = a;
void flip(Tri *A, int a) {
 auto [B, b] = A->e[a]; /* flip edge between A,B */
 if (!B || !A->check(B->p[b])) return;
 Tri *X = new (it++) Tri(A->p[R(a)], B->p[b], A->p[a]);
Tri *Y = new (it++) Tri(B->p[R(b)], A->p[a], B->p[b]);
 link(E(X, 0), E(Y, 0));
 link(E(X, 1), A->e[L(a)]); link(E(X, 2), B->e[R(b)]); link(E(Y, 1), B->e[L(b)]); link(E(Y, 2), A->e[R(a)]);
 A\rightarrow ch = B\rightarrow ch = \{X, Y, nullptr\};
 flip(X, 1); flip(X, 2); flip(Y, 1); flip(Y, 2);
void add_point(int p) {
```

```
Tri *r = root;
while (r->has_chd()) for (Tri *c: r->ch)
 if (c && c->contains(p)) { r = c; break; }
 array<Tri*, 3> t; /* split into 3 triangles */
F3 t[i] = new (it++) Tri(r->p[i], r->p[R(i)], p);
F3 link(E(t[i], 0), E(t[R(i)], 1));
F3 link(E(t[i], 2), r->e[L(i)]);
r->ch = t
F3 flip(t[i], 2);
auto build(const vector<P> &p) {
it = pool; int n = (int)p.size();
vector<int> ord(n); iota(all(ord), 0);
shuffle(all(ord), mt19937(114514));
root = new (it++) Tri(n, n + 1, n + 2);
copy_n(p.data(), n, v); v[n++] = P(-C, -C);
 v[n++] = P(C * 2, -C); v[n++] = P(-C, C * 2);
for (int i : ord) add_point(i);
vector<array<int, 3>> res;
for (Tri *now = pool; now != it; now++)
 if (!now->has_chd()) res.push_back(now->p);
return res;
```

# 6.27 Build Voronoi [94f000]

```
void build_voronoi_cells(auto &&p, auto &&res) {
  vector<vector<int>> adj(p.size());
  for (auto f: res) F3 {
    int a = f[i], b = f[R(i)];
    if (a >= p.size() || b >= p.size()) continue;
    adj[a].emplace_back(b);
}

// use `adj` and `p` and HPI to build cells
for (size_t i = 0; i < p.size(); i++) {
    vector<Line> ls = frame; // the frame
    for (int j : adj[i]) {
        P m = p[i] + p[j], d = rot90(p[j] - p[i]);
        assert (norm(d) != 0);
        ls.emplace_back(m, m + d); // doubled coordinate
    } // HPI(ls)
}
```

# 6.28 kd Tree (Nearest Point) [dbade8]

```
struct KDTree {
struct Node {
  int x, y, x1, y1, x2, y2, id, f; Node *L, *R;
 } tree[maxn], *root;
lld dis2(int x1, int y1, int x2, int y2) {
 lld dx = x1 - x2, dy = y1 - y2;
  return dx * dx + dy * dy;
static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
void init(vector<pair<int,int>> &ip) {
  const int n = ip.size();
 for (int i = 0; i < n; i++) {</pre>
   tree[i].id = i;
   tree[i].x = ip[i].first;
   tree[i].y = ip[i].second;
 }
 root = build(0, n-1, 0);
Node* build(int L, int R, int d) {
   if (L>R) return nullptr; int M = (L+R)/2;
 nth_element(tree+L,tree+M,tree+R+1,d%2?cmpy:cmpx);
 Node &o = tree[M]; o.f = d \% 2;
 o.x1 = o.x2 = o.x; o.y1 = o.y1 = o.y;
 o.L = build(L, M-1, d+1); o.R = build(M+1, R, d+1);
 for (Node *s: {o.L, o.R}) if (s) {
  o.x1 = min(o.x1, s->x1); o.x2 = max(o.x2, s->x2);
o.y1 = min(o.y1, s->y1); o.y2 = max(o.y2, s->y2);
  }
 return tree+M;
bool touch(int x, int y, lld d2, Node *r){
 lld d = sqrt(d2)+1;
return x >= r->x1 - d && x <= r->x2 + d &&
         y >= r - y_1 - d \&\& y <= r - y_2 + d;
using P = pair<lld, int>;
```

void dfs(int x, int y, P &mn, Node \*r) {

```
if (!r || !touch(x, y, mn.first, r)) return;
mn = min(mn, P(dis2(r->x, r->y, x, y), r->id));
if (r->f == 1 ? y < r->y : x < r->x)
    dfs(x, y, mn, r->L), dfs(x, y, mn, r->R);
else
    dfs(x, y, mn, r->R), dfs(x, y, mn, r->L);
}
int query(int x, int y) {
    P mn(INF, -1); dfs(x, y, mn, root);
    return mn.second;
}
} tree;
```

# 6.29 kd Closest Pair (3D ver.) [84d9eb]

```
llf solve(vector<P> v) {
shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,</pre>
 unordered_map<lld, int>>> m;
 llf d = dis(v[0], v[1]);
 auto Idx = [&d] (llf x) -> lld {
  return round(x * 2 / d) + 0.1; };
 auto rebuild_m = [&m, &v, &Idx](int k) {
 m.clear();
for (int i = 0; i < k; ++i)</pre>
   m[Idx(v[i].x)][Idx(v[i].y)]
    [Idx(v[i].z)] = i;
 }; rebuild_m(2);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
     kz = Idx(v[i].z); bool found = false;
  for (int dx = -2; dx \le 2; ++dx) {
   const lld nx = dx + kx;
   if (m.find(nx) == m.end()) continue;
   auto& mm = m[nx];
   for (int dy = -2; dy <= 2; ++dy) {
    const lld ny = dy + ky;
    if (mm.find(ny) == mm.end()) continue;
    auto& mmm = mm[ny];
    for (int dz = -2; dz <= 2; ++dz) {
  const lld nz = dz + kz;</pre>
     if (mmm.find(nz) == mmm.end()) continue;
     const int p = mmm[nz];
     if (dis(v[p], v[i]) < d) {</pre>
      d = dis(v[p], v[i]);
      found = true;
  if (found) rebuild_m(i + 1);
 else m[kx][ky][kz] = i;
 return d:
```

#### 6.30 Simulated Annealing [4e0fe5]

```
llf anneal() {
mt19937 rnd_engine(seed);
uniform_real_distribution<llf> rnd(0, 1);
const llf dT = 0.001;
  / Argument p
llf S_cur = calc(p), S_best = S_cur;
for (llf T = 2000; T > EPS; T -= dT) {
 // Modify p to p_prime
 const llf S_prime = calc(p_prime);
 const llf delta_c = S_prime - S_cur;
 llf prob = min((llf)1, exp(-delta_c / T));
 if (rnd(rnd_engine) <= prob)</pre>
  S_cur = S_prime, p = p_prime;
 if (S_prime < S_best) // find min</pre>
  S_best = S_prime, p_best = p_prime;
}
return S_best;
```

# 6.31 Triangle Centers [adb146]

```
0 = ... // see min circle cover
G = (A + B + C) / 3;
H = G * 3 - 0 * 2; // orthogonal center
llf a = abs(B - C), b = abs(A - C), c = abs(A - B);
I = (a * A + b * B + c * C) / (a + b + c);
// FermatPoint: minimizes sum of distance
```

```
// if max. angle >= 120 deg then vertex
// otherwise, make eq. triangle AB'C, CA'B, BC'A
// line AA', BB', CC' intersects at P
```

# Stringology 7.1 Hash [ce7fad]

```
template <int P = 127, int Q = 1051762951>
class Hash {
vector<int> h, p;
public:
Hash(const auto &s) : h(s.size()+1), p(s.size()+1) {
  for (size_t i = 0; i < s.size(); ++i)</pre>
  h[i + 1] = add(mul(h[i], P), s[i]);
  generate(all(p), [x = 1, y = 1, this]() mutable {
  return y = x, x = mul(x, P), y; });
int query(int l, int r) const { // 1-base (l, r]
 return sub(h[r], mul(h[l], p[r - l]));
};
```

```
Suffix Array [216598]
auto sais(const auto &s) {
 const int n = (int)s.size(), z = ranges::max(s) + 1;
 vector<int> c(z); for (int x : s) ++c[x];
 partial_sum(all(c), begin(c));
 vector<int> sa(n); auto I = ranges::iota_view(0, n);
 if (ranges::max(c) <= 1) {
  for (int i : I) sa[--c[s[i]]] = i;
  return sa;
 vector<bool> t(n); t[n - 1] = true;
 for (int i = n - 2; i >= 0; --i)
  t[i] = (s[i]==s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 auto is_lms = ranges::views::filter([&t](int x) {
  return x && t[x] && !t[x - 1]; });
 const auto induce = [&] {
  for (auto x = c; int y : sa)
   if (y--) if (!t[y]) sa[x[s[y] - 1]++] = y;
  for (auto x = c; int y : sa | views::reverse)
   if (y--) if (t[y]) sa[--x[s[y]]] = y;
 };
 vector<int> lms, q(n); lms.reserve(n);
for (auto x = c; int i : I | is_lms) {
  q[i] = int(lms.size());
  lms.push_back(sa[--x[s[i]]] = i);
 induce(); vector<int> ns(lms.size());
 for (int j = -1, nz = 0; int i : sa | is_lms) {
  if (j >= 0) {
   int len = min({n - i, n - j, lms[q[i] + 1] - i});
   ns[q[i]] = nz += lexicographical_compare(
     begin(s) + j, begin(s) + j + len,
begin(s) + i, begin(s) + i + len);
  j = i;
 }
 ranges::fill(sa, 0); auto nsa = sais(ns);
 for (auto x = c; int y : nsa | views::reverse)
  y = lms[y], sa[--x[s[y]]] = y;
 return induce(), sa;
// sa[i]: sa[i]-th suffix is the
// i-th lexigraphically smallest suffix.
// hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
struct Suffix {
 int n; vector<int> sa, hi, rev;
 Suffix(const string &s) : n(int(s.size())),
  hi(n), rev(n) {
  vector < int > _s(n + 1); _s[n] = 0;
  copy(all(s), begin(_s)); // s shouldn't contain 0
  sa = sais(_s); sa.erase(sa.begin());
  for (int i = 0; i < n; i++) rev[sa[i]] = i;</pre>
  for (int i = 0, h = 0; i < n; ++i) {
   if (!rev[i]) { h = 0; continue; }
   for (int j = sa[rev[i] - 1]; i + h < n</pre>
     && s[i + h] == s[j + h]; ++h);
   hi[rev[i]] = h ? h-- : 0;
};
```

### **7.3** Ex SAM [58374b]

```
struct exSAM {
 int len[maxn * 2], link[maxn * 2]; // maxlen, suflink
 int next[maxn * 2][maxc], tot; // [0, tot), root = 0
 int ord[maxn * 2]; // topo. order (sort by length)
int cnt[maxn * 2]; // occurence
 int newnode() {
  fill_n(next[tot], maxc, 0);
  return len[tot] = cnt[tot] = link[tot] = 0, tot++;
 void init() { tot = 0, newnode(), link[0] = -1; }
 int insertSAM(int last, int c) {
  int cur = next[last][c];
  len[cur] = len[last] + 1;
  int p = link[last];
  while (p != -1 && !next[p][c])
   next[p][c] = cur, p = link[p];
  if (p == -1) return link[cur] = 0, cur;
  int q = next[p][c];
  if (len[p] + 1 == len[q]) return link[cur] = q, cur;
  int clone = newnode();
  for (int i = 0; i < maxc; ++i)</pre>
   next[clone][i] = len[next[q][i]] ? next[q][i] : 0;
  len[clone] = len[p] + 1;
  while (p != -1 && next[p][c] == q)
   next[p][c] = clone, p = link[p];
  link[link[cur] = clone] = link[q];
  link[q] = clone;
  return cur;
 void insert(const string &s) {
  int cur = 0;
  for (char ch : s) {
   int &nxt = next[cur][int(ch - 'a')];
   if (!nxt) nxt = newnode();
   cnt[cur = nxt] += 1;
 void build() {
  queue<int> q; q.push(0);
  while (!q.empty()) {
   int cur = q.front(); q.pop();
   for (int i = 0; i < maxc; ++i)</pre>
    if (next[cur][i]) q.push(insertSAM(cur, i));
  vector<int> lc(tot);
  for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
  partial_sum(all(lc), lc.begin());
  for (int i = 1; i < tot; ++i) ord[--lc[len[i]]] = i;</pre>
 void solve() {
  for (int i = tot - 2; i >= 0; --i)
   cnt[link[ord[i]]] += cnt[ord[i]];
};
7.4 Z value [6a7fd0]
vector<int> Zalgo(const string &s) {
```

```
vector<int> z(s.size(), s.size());
for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
 int j = clamp(r - i, 0, z[i - l]);
 for (; i + j < z[0] and s[i + j] == s[j]; ++j);
if (i + (z[i] = j) > r) r = i + z[l = i];
return z;
```

# 7.5 Manacher [c938a9]

```
vector<int> manacher(const string &S) {
 const int n = (int)S.size(), m = n * 2 + 1;
 vector<int> z(m);
 string t = "."; for (char c: S) t += c, t += '.';
 for (int i = 1, l = 0, r = 0; i < m; ++i) {
 z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
  while (i - z[i] >= 0 \&\& i + z[i] < m) {
   if (t[i - z[i]] == t[i + z[i]]) ++z[i];
  else break:
 if (i + z[i] > r) r = i + z[i], l = i;
}
 return z; // the palindrome lengths are z[i] - 1
```

```
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/* for (int i = 1; i + 1 < m; ++i) {
  int l = (i - z[i] + 2) / 2, r = (i + z[i]) / 2;
  if (l != r) // [l, r) is maximal palindrome
7.6 Lyndon Factorization [d22cc9]
// partition s = w[0] + w[1] + ... + w[k-1],
// w[0] >= w[1] >= ... >= w[k-1]
// each w[i] strictly smaller than all its suffix
void duval(const auto &s, auto &&report) {
 for (int n = (int)s.size(), i = 0, j, k; i < n; ) {</pre>
  for (j = i + 1, k = i; j < n && s[k] <= s[j]; j++)
k = (s[k] < s[j] ? i : k + 1);</pre>
  // if (i < n / 2 && j >= n / 2) {
  // for min cyclic shift, call duval(s + s)
  // then here s.substr(i, n / 2) is min cyclic shift
  for (; i <= k; i += j - k)</pre>
   report(i, j - k); // s.substr(l, len)
} // tested @ luogu 6114, 1368 & UVA 719
7.7 Main Lorentz [615b8f]
vector<pair<int, int>> rep[kN]; // 0-base [l, r]
void main_lorentz(const string &s, int sft = 0) {
 const int n = s.size();
 if (n == 1) return;
 const int nu = n / 2, nv = n - nu;
 const string u = s.substr(0, nu), v = s.substr(nu);
   ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend());
 main_lorentz(u, sft), main_lorentz(v, sft + nu);
 const auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u),
    z3 = Zalgo(ru + '#' + rv), z4 = Zalgo(v);
 auto get_z = [](const vector<int> &z, int i) {
  return (0 <= i and i < (int)z.size()) ? z[i] : 0; };</pre>
 auto add_rep = [&](bool left, int c, int l, int k1,
    int k2) {
  const int L = max(1, l - k2), R = min(l - left, k1);
  if (L > R) return;
  if (left) rep[l].emplace_back(sft + c - R, sft + c -
    L);
  else rep[l].emplace_back(sft + c - R - l + 1, sft + c
      - L - l + 1);
 for (int cntr = 0; cntr < n; cntr++) {</pre>
  int l, k1, k2;
  if (cntr < nu) {</pre>
   l = nu - cntr;
   k1 = get_z(z1, nu - cntr);
   k2 = get_z(z2, nv + 1 + cntr);
  } else {
   l = cntr - nu + 1;
   k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
   k2 = get_z(z4, (cntr - nu) + 1);
  if (k1 + k2 >= l)
   add_rep(cntr < nu, cntr, l, k1, k2);</pre>
}
7.8 BWT [5a9b3a]
vector<int> v[SIGMA];
void BWT(char *ori, char *res) {
// make ori -> ori + ori
 // then build suffix array
void iBWT(char *ori, char *res) {
 for (int i = 0; i < SIGMA; i++) v[i].clear();</pre>
 const int len = strlen(ori);
```

```
for (int i = 0; i < len; i++)
 v[ori[i] - 'a'].push_back(i);
vector<int> a;
for (int i = 0, ptr = 0; i < SIGMA; i++)</pre>
 for (int j : v[i]) {
  a.push_back(j);
  ori[ptr++] = 'a' + i;
for (int i = 0, ptr = 0; i < len; i++) {</pre>
res[i] = ori[a[ptr]];
 ptr = a[ptr];
res[len] = 0;
```

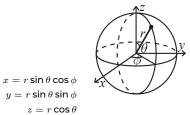
# 7.9 Palindromic Tree [0673ee]

```
struct PalindromicTree {
 struct node {
  int nxt[26], f, len; // num = depth of fail link
                  // = #pal_suffix of this node
  int cnt, num;
  node(int l = 0) : nxt{}, f(0), len(l), cnt(0), num(0)
 vector<node> st; vector<char> s; int last, n;
 void init() {
  st.clear(); s.clear();
  last = 1; n = 0;
  st.push_back(0); st.push_back(-1);
  st[0].f = 1; s.push_back(-1);
 int getFail(int x) {
  while (s[n - st[x].len - 1] != s[n]) x = st[x].f;
  return x;
 void add(int c) {
  s.push_back(c -= 'a'); ++n;
  int cur = getFail(last);
  if (!st[cur].nxt[c]) {
   int now = st.size();
   st.push_back(st[cur].len + 2);
   st[now].f = st[getFail(st[cur].f)].nxt[c];
   st[cur].nxt[c] = now;
   st[now].num = st[st[now].f].num + 1;
  last = st[cur].nxt[c]; ++st[last].cnt;
 void dpcnt() { // cnt = #occurence in whole str
  for (int i = st.size() - 1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
 int size() { return st.size() - 2; }
} pt;
string s; cin >> s; pt.init();
for (int i = 0; i < SZ(s); i++) {</pre>
 int prvsz = pt.size(); pt.add(s[i]);
 if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
  // pal @ [l,r]: s.substr(l, r-l+1)
}
} */
```

#### 8 Misc

#### 8.1 Theorems

#### **Spherical Coordinate**



$$\begin{split} r &= \sqrt{x^2 + y^2 + z^2} \\ \theta &= \mathsf{acos}(z/\sqrt{x^2 + y^2 + z^2}) \\ \phi &= \mathsf{atan2}(y,x) \end{split}$$

#### Sherman-Morrison formula

$$(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1+v^{\mathsf{T}}A^{-1}u}$$

# Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} =$ d(i),  $L_{ij} = -c$  where c is the number of edge (i, j) in  $\tilde{G}$ .

- The number of undirected spanning in G is  $\det(\tilde{L}_{11})$ .
- The number of directed spanning tree rooted at r in G is  $det(\tilde{L}_{rr})$ .

#### Random Walk on Graph

Let P be the transition matrix of a strongly connected directed graph,  $\sum_j P_{i,j}=1$ . Let  $F_{i,j}$  be the expected time to reach j from i. Let  $g_i$  be the expected time from i to i, G = diag(g) and J be a matrix all of 1, i.e.  $J_{i,j} = 1$ . Then, F = J - G + PF

First solve G: let  $\pi P=\pi$  be a stationary distribution. Then  $\pi_i g_i=1$ . The rank of I-P is n-1, so we first solve a special solution X such that (I-P)X=J-G and adjust X to F by  $F_{i,j}=X_{i,j}-X_{j,j}$ .

#### Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in \mathit{E}$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

#### Cayley's Formula

- · Given a degree sequence  $d_1, d_2, \ldots, d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees. • Let  $T_{n,k}$  be the number of labeled forests on n vertices with k com-
- ponents, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$  .

#### Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1 + d_2 + \ldots + d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all  $1 \leq k \leq n$ .

#### Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

#### Euler's planar graph formula

V - E + F = C + 1.  $E \le 3V - 6$  (when  $V \ge 3$ )

#### Pick's theorem

For simple polygon, when points are all integer, we have  ${\it A}$ #{lattice points in the interior} +  $\frac{\#\{\text{lattice points on the boundary}\}}{2} - 1$ 

#### Matroid Intersection

Given matroids  $M_1=(G,I_1), M_2=(G,I_2)$ , find maximum  $S\in I_1\cap I_2$ . For each iteration, build the directed graph and find a shortest path from s to t. •  $s \rightarrow x : S \sqcup \{x\} \in I_1$ •  $x \to t : S \sqcup \{x\} \in I_2$ 

•  $x \to t$  .  $S \sqcup \{x\} \in I_2$ •  $y \to x$  :  $S \setminus \{y\} \sqcup \{x\} \in I_1$  (y is in the unique circuit of  $S \sqcup \{x\}$ ) •  $x \to y$  :  $S \setminus \{y\} \sqcup \{x\} \in I_2$  (y is in the unique circuit of  $S \sqcup \{x\}$ ) Alternate the path, and |S| will increase by 1. Let R = min(rank( $I_1$ ), rank( $I_2$ ), N = |G|. In each iteration,  $|E| = O(R^{-1})$ For weighted case, assign weight -w(x) and w(x) to  $x \in S$  and  $x \notin S$ , resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1. **8.2 Weight Matroid Intersection** [d00ee8]

```
struct Matroid {
 Matroid(bitset<N>); // init from an independent set
 bool can_add(int); // check if break independence
 Matroid remove(int); // removing from the set
auto matroid_intersection(const vector<int> &w) {
 const int n = (int)w.size(); bitset<N> S;
 for (int sz = 1; sz <= n; sz++) {</pre>
  Matroid M1(S), M2(S); vector<vector<pii>> e(n + 2);
  for (int j = 0; j < n; j++) if (!S[j]) {</pre>
   if (M1.can_add(j)) e[n].eb(j, -w[j]);
   if (M2.can_add(j)) e[j].eb(n + 1, 0);
  for (int i = 0; i < n; i++) if (S[i]) {</pre>
   Matroid T1 = M1.remove(i), T2 = M2.remove(i);
   for (int j = 0; j < n; j++) if (!S[j]) {
  if (T1.can_add(j)) e[i].eb(j, -w[j]);
}</pre>
    if (T2.can_add(j)) e[j].eb(i, w[i]);
  } // maybe implicit build graph for more speed
  vector<pii> d(n + 2, \{INF, 0\}); d[n] = \{0, 0\};
  vector<int> prv(n + 2, -1);
  // change to SPFA for more speed, if necessary
  for (int upd = 1; upd--; )
   for (int u = 0; u < n + 2; u++)
    for (auto [v, c] : e[u]) {
     pii x(d[u].first + c, d[u].second + 1);
if (x < d[v]) d[v] = x, prv[v] = u, upd = 1;</pre>
  if (d[n + 1].first >= INF) break;
  for (int x = prv[n+1]; x!=n; x = prv[x]) S.flip(x);
  // S is the max-weighted independent set w/ size sz
 return S:
} // from Nacl
```

#### Stable Marriage

```
1: Initialize m \in M and w \in W to free
2: while \exists free man m who has a woman w to propose to do
       w \leftarrow \text{first woman on } m \text{'s list to whom } m \text{ has not yet proposed}
      if \exists some pair (m', w) then
          if w prefers m to m' then
              m' \leftarrow free
              (m,w) \gets \mathsf{engaged}
          end if
```

```
else
            (m,w) \gets \mathsf{engaged}
10:
        end if
11:
12: end while
```

### 8.4 Bitset LCS [330ab1]

```
cin >> n >> m;
for (int i = 1, x; i <= n; ++i)</pre>
 cin >> x, p[x].set(i);
for (int i = 1, x; i <= m; i++) {
 cin >> x, (g = f) |= p[x];
 f.shiftLeftByOne(), f.set(0);
 ((f = g - f) ^= g) \&= g;
cout << f.count() << '\n';</pre>
```

### 8.5 Prefix Substring LCS [7d8faf]

```
void all_lcs(string S, string T) { // 0-base
  vector<size_t> h(T.size()); iota(all(h), 1);
  for (size_t a = 0; a < S.size(); ++a) {</pre>
   for (size_t c = 0, v = 0; c < T.size(); ++c)</pre>
    if (S[a] == T[c] || h[c] < v) swap(h[c], v);</pre>
   // here, LCS(s[0, a], t[b, c]) =
   // c - b + 1 - sum([h[i] > b] | i <= c)
} // test @ yosupo judge
```

# 8.6 Convex 1D/1D DP [6e0124]

```
struct segment {
 int i, l, r;
 segment() {}
 segment(int a, int b, int c): i(a), l(b), r(c) {}
void solve()
 auto f = [](int l, int r){return dp[l] + w(l+1, r);}
 dp[0] = 0;
 deque<segment> dq; dq.push_back(segment(0, 1, n));
for (int i = 1; i <= n; ++i) {</pre>
  dp[i] = f(dq.front().i, i);
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
   f(i, dq.back().l)<f(dq.back().i, dq.back().l))</pre>
    dq.pop_back();
  if (dq.size()) {
   int d = 1 << 20, c = dq.back().l;</pre>
   while (d >>= 1) if (c + d <= dq.back().r)</pre>
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
  if (seg.l <= n) dq.push_back(seg);</pre>
}
```

#### ConvexHull Optimization [b4318e]

```
mutable lld a, b, p;
 bool operator<(const L &r) const {</pre>
  return a < r.a; /* here */
 bool operator<(lld x) const { return p < x; }</pre>
lld Div(lld a, lld b) {
return a / b - ((a ^ b) < 0 && a % b); }
struct DynamicHull : multiset<L, less<>>> {
 static const lld kInf = 1e18;
 bool Isect(iterator x, iterator y) {
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a)
   x->p = x->b > y->b ? kInf : -kInf; /* here */
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(lld a, lld b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
  while ((y = x) != begin() && (--x)->p >= y->p)
   Isect(x, erase(y));
 lld Query(lld x) { // default chmax
 auto l = *lower_bound(x); // to chmin:
return l.a * x + l.b; // modify the 2 "<>"
```

```
];
```

```
8.8 Min Plus Convolution [464dcd]

// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]

vector<int> min_plus_convolution(auto &a, auto &b) {
    const int n = (int)a.size(), m = (int)b.size();
    vector<int> c(n + m - 1, numeric_limits<int>::max());
    auto dc = [&](auto Y, int l, int r, int jl, int jr) {
        if (l > r) return;
        int mid = (l + r) / 2, from = -1, &best = c[mid];
        for (int j = jl; j <= jr; j++)
        if (int i = mid - j; i >= 0 && i < n)
              if (best > a[i]+b[j]) best = a[i]+b[j], from = j;
        Y(Y, l, mid-1, jl, from); Y(Y, mid+1, r, from, jr);
    };
    return dc(dc, 0, n-1+m-1, 0, m-1), c;
}
```

# 8.9 De-Bruijn [c0a223]

```
vector<int> de_bruijn(int k, int n) {
  // return cyclic string of len k^n s.t. every string
  // of len n using k char appears as a substring.
  vector<int> aux(n + 1), res;
  auto db = [&](auto self, int t, int p) -> void {
    if (t <= n)
      for (int i = aux[t - p]; i < k; ++i, p = t)
        aux[t] = i, self(self, t + 1, p);
    else if (n % p == 0) for (int i = 1; i <= p; ++i)
        res.push_back(aux[i]);
    }; db(db, 1, 1);
    return res;
}</pre>
```

# 8.10 Josephus Problem [f4494f]

```
int f(int n, int m) { // n people kill m for each turn
int s = 0;
for (int i = 2; i <= n; i++) s = (s + m) % i;
return s;
}
int kth(int n, int m, int k){ // died at kth
if (m == 1) return n-1;
for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
return k;
}
```

#### 8.11 N Queens Problem [31f83e]

```
void solve(VI &ret, int n) { // no sol when n=2,3
if (n % 6 == 2) {
    for (int i = 2; i <= n; i += 2) ret.push_back(i);
    ret.push_back(3); ret.push_back(1);
    for (int i = 7; i <= n; i += 2) ret.push_back(i);
    ret.push_back(5);
} else if (n % 6 == 3) {
    for (int i = 4; i <= n; i += 2) ret.push_back(i);
    ret.push_back(2);
    for (int i = 5; i <= n; i += 2) ret.push_back(i);
    ret.push_back(1); ret.push_back(3);
} else {
    for (int i = 2; i <= n; i += 2) ret.push_back(i);
    for (int i = 1; i <= n; i += 2) ret.push_back(i);
}</pre>
```

#### 8.12 Tree Knapsack [f42766]

```
vector<int> G[N]; int dp[N][K]; pair<int,int> obj[N];
void dfs(int u, int mx) {
  for (int s : G[u]) {
    auto [w, v] = obj[s];
    if (mx < w) continue;
    for (int i = 0; i <= mx - w; i++)
        dp[s][i] = dp[u][i];
    dfs(s, mx - w);
  for (int i = w; i <= mx; i++)
        dp[u][i] = max(dp[u][i], dp[s][i - w] + v);
    }
}</pre>
```

#### 8.13 Manhattan MST [1008bc]

```
vector<array<int, 3>> manhattanMST(vector<P> ps) {
  vector<int> id(ps.size()); iota(all(id), 0);
  vector<array<int, 3>> edges;
```

```
for (int k = 0; k < 4; k++) {
    sort(all(id), [&](int i, int j) {
        return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y; });
    map<int, int> sweep;
    for (int i : id) {
        for (auto it = sweep.lower_bound(-ps[i].y);
            it != sweep.end(); sweep.erase(it++)) {
            if (P d = ps[i] - ps[it->second]; d.y > d.x) break;
            else edges.push_back({d.y + d.x, i, it->second});
        }
        sweep[-ps[i].y] = i;
    }
    for (P &p : ps)
        if (k & 1) p.x = -p.x;
        else swap(p.x, p.y);
    }
    return edges; // [{w, i, j}, ...]
} // test @ yosupo judge
```

## 8.14 Binary Search On Fraction [765c5a]

```
struct Q {
ll p, q;
 Q go(Q b, ll d) { return {p + b.p*d, q + b.q*d}; }
bool pred(0);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
 Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
 assert(pred(hi));
 bool dir = 1, L = 1, H = 1;
 for (; L || H; dir = !dir) {
  ll len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
   if (Q mid = hi.go(lo, len + step);
     mid.p > N || mid.q > N || dir ^ pred(mid))
    t++;
   else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
 return dir ? hi : lo;
```

#### 8.15 Barrett Reduction [d44617]

```
struct FastMod {
  using Big = __uint128_t; llu b, m;
  FastMod(llu b) : b(b), m(-1ULL / b) {}
  llu reduce(llu a) { // a % b
   llu r = a - (llu)((Big(m) * a) >> 64) * b;
  return r >= b ? r - b : r;
  }
};
```

#### 8.16 Montgomery [1f2b4a]

```
struct Mont { // Montegomery multiplication
 constexpr static int W = 64;
 llu mod, R1Mod, R2Mod, NPrime;
 void set_mod(llu _mod) {
 mod = _mod; assert(mod & 1);
  const u128 B = u128(1) << W, R = B % mod;</pre>
  llu xinv = 1;
  for (int i = 1; i < W; i++) // Hensel lifting</pre>
   if ((xinv * mod) >> i & 1) xinv |= 1ULL << i;</pre>
  assert(xinv * mod == 1);
  R1Mod = static_cast<llu>(R);
  R2Mod = static_cast<llu>(R * R % mod);
 NPrime = static_cast<llu>(B - xinv);
 llu redc(llu a, llu b) const {
  auto T = static_cast<u128>(a) * b;
  u128 m = static_cast<llu>(T) * NPrime;
  T += m * mod; T >>= W;
  return static_cast<llu>(T >= mod ? T - mod : T);
 llu from(llu x) const {
 assert(x < mod);</pre>
 return redc(x, R2Mod);
 llu get(llu a) const { return redc(a, 1); }
llu one() const { return R1Mod; }
```

```
|} mont;
|// a * b % mod == get(redc(from(a), from(b)))
```