Contents		6.19 Circle Cover
1 Basic 1.1 vimrc 1.2 Debug Macro 1.3 Increase Stack 1.4 Pragma Optimization 1.5 IO Optimization	1 1 2	7 Stringology 22 7.1 Hash 22 7.2 Suffix Array 22 7.3 Suffix Automaton 22 7.4 KMP 23 7.5 Z value 23 7.6 Manacher 23 7.7 Lexico Smallest Rotation 23
2 Data Structure 2.1 Dark Magic 2.2 Link-Cut Tree 2.3 LiChao Segment Tree 2.4 Treap 2.5 Linear Basis 2.6 Binary Search On Segment Tree	2 3 3 3	7.8 Main Lorentz 23 7.9 BWT 23 7.10 Palindromic Tree 24 8 Misc 24 8.1 Theorems 24 8.1.1 Sherman-Morrison formula 24
3 Craph 3.1 2-SAT (SCC) 3.2 BCC 3.3 Round Square Tree 3.4 Centroid Decomposition 3.5 Directed Minimum Spanning Tree 3.6 Dominator Tree 3.7 Edge Coloring 3.8 Lowbit Decomposition 3.9 Manhattan Minimum Spanning Tree 3.10 MaxClique 3.11 MaxClique 3.12 Minimum Mean Cycle 3.13 Mo's Algorithm on Tree 3.14 Virtual Tree	4 4 4 5 5 5 5 6 6 6 7 7 7 8	8.1.2 Kirchhoff's Theorem 24 8.1.3 Tutte's Matrix 24 8.1.4 Cayley's Formula 24 8.1.5 Erdős-Gallai theorem 24 8.1.6 Havel-Hakimi algorithm 24 8.1.7 Euler's planar graph formula 24 8.1.8 Pick's theorem 24 8.1.9 Matroid Intersection 24 8.2 Bitset LCS 24 8.3 Prefix Substring LCS 24 8.4 Convex 1D/1D DP 24 8.5 Convex Hull Optimization 25 8.6 Josephus Problem 25 8.7 Tree Knapsack 25 8.8 N Queens Problem 25 8.9 Stable Marriage 25
4 Matching & Flow 4.1 Bipartite Matching 4.2 Dijkstra Cost Flow 4.3 Dinic 4.4 Construct VC	8 8	8.10 Binary Search On Fraction
4.5 Flow Models 4.6 General Graph Matching 4.7 Global Min-Cut 4.8 GomoryHu Tree 4.9 Kuhn Munkres 4.10 Minimum Cost Circulation 4.11 Minimum Cost Maximum Flow 4.12 Maximum Weight Graph Matching	9 10 10 10 10	<pre>1.1 vimrc se is nu ru et tgc sc hls cin cino+=j1 sw=4 sts=4 bs=2 mouse=a "encoding=utf-8 ls=2 syn on colo desert filetype indent on inoremap {<cr> {<cr>}<esc>0 map <f8> <esc>:w<cr>:!g++ "%" -o "%<" -std=c++17 -</cr></esc></f8></esc></cr></cr></pre>
5 Math 5.1 Common Bounds 5.1.1 Partition function 5.1.2 Divisor function 5.1.3 Factorial 5.1.4 Binom Coef 5.2 Strling Number 5.2.1 First Kind 5.2.2 Second Kind 5.3 ax+by=gcd	13 13 13 13 13 13 13	DCKISEKI -Wall -Wextra -Wshadow -Wfatal-errors - Wconversion -fsanitize=address, undefined -g && echo success <cr> map <f9> <esc>:w<cr>:!g++ "%" -o "%<" -02 -std=c++17 && echo success<cr> map <f10> <esc>:!./"%<"<cr> 1.2 Debug Macro</cr></esc></f10></cr></cr></esc></f9></cr>
5.5 ax+by-gcd 5.4 Berlekamp Massey 5.5 Charateristic Polynomial 5.6 Chinese Remainder 5.7 De-Bruijn 5.8 DiscreteLog 5.9 Extended Euler 5.10 Extended FloorSum 5.11 Fast Fourier Transform 5.12 FloorSum 5.13 FWT 5.14 Miller Rabin 5.15 NTT 5.16 Partition Number 5.17 Pi Count (Linear Sieve) 5.18 Pollard Rho 5.19 Polynomial Operations 5.20 Quadratic residue 5.21 Simplex 5.22 Simplex Construction	13 13 14 14 14 14 15 15 15 15 16 16 16 17	#ifdef KISEKI #define safe cerr< <pretty_function\< td=""></pretty_function\<>
6 Geometry 6.1 Basic Geometry 6.2 Segment & Line Intersection 6.3 2D Convex Hull 6.4 3D Convex Hull 6.5 2D Farthest Pair 6.6 kD Closest Pair (3D ver.) 6.7 Simulated Annealing 6.8 Half Plane Intersection 6.9 Minkowski Sum	18 18 18 18 18 19	<pre>cerr << (f++ ? ", " : "") << *L; cerr << "]\e[0m\n"; } #else #define safe ((void)0) #define debug() ((void)0) #define orange() ((void)0) #endif</pre>
6.10 Circle Class 6.11 Intersection of line and Circle 6.12 Intersection of Polygon and Circle 6.13 Point & Hulls Tangent 6.14 Convex Hulls Tangent 6.15 Tangent line of Two Circle 6.16 Minimum Covering Circle 6.17 KDTree (Nearest Point) 6.18 Rotating Sweep Line	19 19 19 20 20 20 21 21	<pre>1.3 Increase Stack const int size = 256 << 20; register long rsp asm("rsp"); char *p = (char*)malloc(size)+size, *bak = (char*)rsp; _asm("movq %0, %%rsp\n"::"r"(p)); // main _asm("movq %0, %%rsp\n"::"r"(bak));</pre>

1.4 Pragma Optimization

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8000)
```

1.5 IO Optimization

```
| static inline int gc() {
    constexpr int B = 1<<20;
    static char buf[B], *p, *q;
    if(p == q &&
        (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
    return EOF;
    return *p++;
}

template < typename T >
    static inline bool gn( T &x ) {
    int c = gc(); T sgn = 1; x = 0;
    while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
    if(c == '-') sgn = -1, c = gc();
    if(c == EOF) return false;
    while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
    return x *= sgn, true;
}
```

2 Data Structure

2.1 Dark Magic

2.2 Link-Cut Tree

```
template <typename Val> class LCT {
private:
struct node {
 int pa, ch[2];
 bool rev;
 Val v, v_prod, v_rprod;
 node() : pa{0}, ch{0, 0}, rev{false}, v{}, v_prod{},
    v_rprod{} {};
vector<node> nodes;
set<pair<int, int>> edges;
bool is_root(int u) const {
 const int p = nodes[u].pa;
 return nodes[p].ch[0] != u and nodes[p].ch[1] != u;
bool is_rch(int u) const {
 return (not is_root(u)) and nodes[nodes[u].pa].ch[1]
    == u;
void down(int u) {
 if (auto &cnode = nodes[u]; cnode.rev) {
  if (cnode.ch[0]) set_rev(cnode.ch[0]);
  if (cnode.ch[1]) set_rev(cnode.ch[1]);
  cnode.rev = false;
 }
void up(int u) {
 auto &cnode = nodes[u];
 cnode.v_prod =
  nodes[cnode.ch[0]].v_prod * cnode.v * nodes[cnode.ch
    [1]].v_prod;
 cnode.v_rprod =
  nodes[cnode.ch[1]].v_rprod * cnode.v * nodes[cnode.
    ch[0]].v_rprod;
}
```

```
void set_rev(int u) {
  swap(nodes[u].ch[0], nodes[u].ch[1]);
  swap(nodes[u].v_prod, nodes[u].v_rprod);
  nodes[u].rev ^= 1;
 void rotate(int u) {
  int f = nodes[u].pa, g = nodes[f].pa, 1 = is_rch(u);
if (nodes[u].ch[1 ^ 1])
   nodes[nodes[u].ch[1 ^ 1]].pa = f;
  if (not is_root(f))
   nodes[g].ch[is_rch(f)] = u;
  nodes[f].ch[1] = nodes[u].ch[1 ^ 1];
  nodes[u].ch[1^{^{\prime}}] = f
  nodes[u].pa = g, nodes[f].pa = u;
  up(f);
 void splay(int u) {
  vector<int> stk = {u};
  while (not is_root(stk.back()))
   stk.push_back(nodes[stk.back()].pa);
  for (; not stk.empty(); stk.pop_back())
   down(stk.back());
  for(int f=nodes[u].pa;!is_root(u);f=nodes[u].pa){
  if(!is_root(f))rotate(is_rch(u)==is_rch(f)?f:u);
   rotate(u);
  up(u);
 void access(int u) {
  int last = 0;
  for (int last = 0; u; last = u, u = nodes[u].pa) {
   splay(u);
   nodes[u].ch[1] = last;
   up(u);
 int find_root(int u) {
  access(u); splay(u);
  int la = 0:
  for (; u; la = u, u = nodes[u].ch[0]) down(u);
  return la;
 void change_root(int u) {
  access(u); splay(u); set_rev(u);
 void link(int x, int y)
  change_root(y); nodes[y].pa = x;
 void split(int x, int y) {
  change_root(x); access(y); splay(y);
 void cut(int x, int y) {
  split(x, y)
  nodes[y].ch[0] = nodes[x].pa = 0;
  up(y);
public:
 LCT(int n = 0) : nodes(n + 1) {}
 int add(const Val &v = {}) {
  nodes.push_back(v);
  return int(nodes.size()) - 2;
 int add(Val &&v) {
 nodes.emplace_back(move(v));
  return int(nodes.size()) - 2;
 void set_val(int u, const Val &v) {
  splay(++u); nodes[u].v = v; up(u);
 Val query(int x, int y) {
  split(++x, ++y);
  return nodes[y].v_prod;
 bool connected(int u, int v) { return find_root(++u)
    == find_root(++v); }
 void add_edge(int u, int v) {
  if (++u > ++v) swap(u, v)
  edges.emplace(u, v); link(u, v);
 void del_edge(int u, int v) {
  auto k = minmax(++u, ++v)
  if (auto it = edges.find(k); it != edges.end()) {
```

```
edges.erase(it); cut(u, v);
                                                                // sz(L) == k
                                                               int getRank(node *o) { // 1-base
 }
                                                               int r = sz(o->lc) + 1;
};
                                                                for (;o->pa != nullptr; o = o->pa)
                                                                if (o->pa->rc == o) r += sz(o->pa->lc) + 1;
      LiChao Segment Tree
struct L {
 int m, k, id;
                                                               #undef sz
 L() : id(-1) \{ \}
L(int a, int b, int c) : m(a), k(b), id(c) {}
int at(int x) { return m * x + k; }
                                                             2.5 Linear Basis
                                                             template <int BITS> struct Basis {
                                                               array<pair<uint64_t, int>, BITS> b;
class LiChao {
private:
                                                               Basis() { b.fill({0, -1});
                                                              void add(uint64_t x, int p) {
  for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
    if (b[i].first == 0) {
 int n; vector<L> nodes;
 static int lc(int x) { return 2 * x + 1; }
 static int rc(int x) { return 2 * x + 2;
 void insert(int 1, int r, int id, L ln) {
                                                                 b[i] = \{x, p\};
  int m = (1 + r) >> 1;
                                                                  return;
  if (nodes[id].id == -1) {
                                                                 } else if (b[i].second > p) {
                                                                  swap(b[i].first, x), swap(b[i].second, p);
  nodes[id] = ln;
   return;
                                                                 x ^= b[i].first;
  bool atLeft = nodes[id].at(1) < ln.at(1);</pre>
                                                               }
  if (nodes[id].at(m) < ln.at(m)) {</pre>
  atLeft ^= 1
                                                              bool ok(uint64_t x, int p) {
  for (int i = 0; i < BITS; ++i)</pre>
   swap(nodes[id], ln);
                                                                if (((x >> i) \& 1) and b[i].second < p)
                                                                 x ^= b[i].first;
  if (r - 1 == 1) return;
  if (atLeft) insert(1, m, lc(id), ln);
                                                                return x == 0;
  else insert(m, r, rc(id), ln);
                                                             };
 int query(int 1, int r, int id, int x) {
                                                             2.6 Binary Search On Segment Tree
 int ret = 0, m = (1 + r) >> 1;
                                                             // find_first = x \rightarrow minimal x s.t. check([a, x))
  if (nodes[id].id != -1)
                                                             // find_last = x \rightarrow maximal x s.t. check([x, b))
   ret = nodes[id].at(x);
  if (r - 1 == 1) return ret;
                                                             template <typename C>
  if (x < m) return max(ret, query(1, m, lc(id), x));</pre>
                                                             int find_first(int 1, const C &check) {
  return max(ret, query(m, r, rc(id), x));
                                                               if (1 >= n) return n + 1;
                                                               1 += sz:
                                                               for (int i = height; i > 0; i--)
public:
                                                               propagate(1 >> i);
LiChao(int n_{-}) : n(n_{-}), nodes(n * 4) \{}
                                                               Monoid sum = identity;
 void insert(L ln) { insert(0, n, 0, ln); }
                                                                while ((1 & 1) == 0) 1 >>= 1;
 int query(int x) { return query(0, n, 0, x); }
                                                                if (check(f(sum, data[1]))) {
                                                                while (1 < sz) {
2.4 Treap
                                                                  propagate(1);
namespace Treap{
 #define sz(x)((x)?((x)-size):0)
                                                                  auto nxt = f(sum, data[1]);
 struct node{
                                                                  if (not check(nxt)) {
  int size;
                                                                   sum = nxt;
  uint32_t pri;
                                                                   1++;
  node *1c, *rc, *pa;
  node():size(0),pri(rand()),lc(0),rc(0),pa(0){}
  void pull() {
                                                                 return 1 + 1 - sz;
  size = 1; pa = nullptr;
   if ( lc ) { size += lc->size; lc->pa = this; }
                                                                sum = f(sum, data[1++]);
                                                               } while ((1 & -1) != 1);
   if ( rc ) { size += rc->size; rc->pa = this; }
                                                               return n + 1;
 node* merge( node* L, node* R )
                                                              template <typename C>
 if ( not L or not R ) return L ? L : R;
                                                             int find_last(int r, const C &check) {
  if ( L->pri > R->pri ) {
                                                              if (r <= 0) return -1;</pre>
   L->rc = merge( L->rc, R ); L->pull();
                                                               r += sz;
   return L;
                                                               for (int i = height; i > 0; i--)
  } else {
                                                                propagate((r - 1) >> i);
   R->lc = merge( L, R->lc ); R->pull();
                                                               Monoid sum = identity;
   return R;
                                                               do {
  }
                                                                while (r > 1 and (r & 1)) r >>= 1;
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                                if (check(f(data[r], sum))) {
  if ( not rt ) L = R = nullptr;
                                                                 while (r < sz) {</pre>
  else if( sz( rt->lc ) + 1 <= k ) {</pre>
                                                                  propagate(r);
                                                                  r = (r << 1) + 1;
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                  auto nxt = f(data[r], sum);
   L->pull();
                                                                  if (not check(nxt)) {
  } else {
                                                                   sum = nxt;
   R = rt:
                                                                   r--;
   split_by_size( rt->lc, k, L, R->lc );
   R->pull();
                                                                 return r - sz;
```

```
sum = f(data[r], sum);
} while ((r & -r) != r);
return -1;
    Graph
3.1 2-SAT (SCC)
class TwoSat{
private:
int n;
vector<vector<int>> rG,G,sccs;
vector<int> ord,idx;
vector<bool> vis,result;
void dfs(int u){
 vis[u]=true
 for(int v:G[u])
  if(!vis[v]) dfs(v);
 ord.push_back(u);
void rdfs(int u){
 vis[u]=false;idx[u]=sccs.size()-1;
 sccs.back().push_back(u);
 for(int v:rG[u])
  if(vis[v])rdfs(v);
public:
void init(int n_){
 G.clear();G.resize(n=n_);
 rG.clear();rG.resize(n);
 sccs.clear();ord.clear();
 idx.resize(n);result.resize(n);
void add_edge(int u,int v){
 G[u].push_back(v);rG[v].push_back(u);
void orr(int x,int y){
 if ((x^y)==1)return;
 add\_edge(x^1,y); \ add\_edge(y^1,x);
bool solve(){
 vis.clear();vis.resize(n);
 for(int i=0;i<n;++i)</pre>
  if(not vis[i])dfs(i);
  reverse(ord.begin(),ord.end());
 for (int u:ord){
  if(!vis[u])continue;
  sccs.push_back(vector<int>());
  rdfs(u);
 for(int i=0;i<n;i+=2)</pre>
  if(idx[i]==idx[i+1])
    return false;
  vector<bool> c(sccs.size());
```

3.2 BCC

return true;

}

} sat2:

}

```
class BCC {
private:
 int n, ecnt;
 vector<vector<pair<int, int>>> g;
 vector<int> dfn, low;
 vector<bool> ap, bridge;
 void dfs(int u, int f) {
  dfn[u] = low[u] = dfn[f] + 1;
  int ch = 0;
  for (auto [v, t] : g[u]) if (v != f) {
  if (dfn[v]) {
    low[u] = min(low[u], dfn[v]);
```

for(size_t i=0;i<sccs.size();++i){</pre>

bool get(int x){return result[x];} int get_id(int x){return idx[x];} int count(){return sccs.size();}

for(auto sij : sccs[i]){ result[sij]=c[i]; c[idx[sij^1]]=!c[i];

```
} else {
    ++ch, dfs(v, u);
    low[u] = min(low[u], low[v]);
    if (low[v] > dfn[u])
bridge[t] = true;
    if (low[v] >= dfn[u])
     ap[u] = true;
  ap[u] &= (ch != 1 or u != f);
public:
 void init(int n_) {
  g.assign(n = n_, vector<pair<int, int>>());
  low.assign(n, ecnt = 0);
  dfn.assign(n, 0);
  ap.assign(n, false);
 void add_edge(int u, int v) {
  g[u].emplace_back(v, ecnt);
  g[v].emplace_back(u, ecnt++);
 void solve() {
  bridge.assign(ecnt, false);
  for (int i = 0; i < n; ++i)</pre>
   if (not dfn[i]) dfs(i, i);
 bool is_ap(int x) { return ap[x]; }
 bool is_bridge(int x) { return bridge[x]; }
```

3.3 Round Square Tree

```
int N, M, cnt;
std::vector<int> G[maxn], T[maxn * 2];
int dfn[maxn], low[maxn], dfc;
int stk[maxn], tp;
void Tarjan(int u) {
 low[u] = dfn[u] = ++dfc;
 stk[++tp] = u;
 for (int v : G[u]) {
  if (!dfn[v]) {
   Tarjan(v);
   low[u] = std::min(low[u], low[v]);
   if (low[v] == dfn[u]) {
    ++cnt:
    for (int x = 0; x != v; --tp) {
     x = stk[tp];
     T[cnt].push_back(x);
     T[x].push_back(cnt);
    T[cnt].push_back(u);
    T[u].push_back(cnt);
  } else
   low[u] = std::min(low[u], dfn[v]);
int main() { // ...
 cnt = N:
 for (int u = 1; u <= N; ++u)
  if (!dfn[u]) Tarjan(u), --tp;
```

3.4 Centroid Decomposition

```
struct Centroid {
 vector<vector<int64_t>> Dist;
 vector<int> Pa, Dep;
 vector<int64_t> Sub, Sub2;
 vector<int> Cnt, Cnt2;
 vector<int> vis, sz, mx, tmp
 void DfsSz(int x) {
 vis[x] = true; sz[x] = 1; mx[x] = 0;
  for (auto [u, w] : g[x]) {
   if (vis[u]) continue;
  DfsSz(u);
   sz[x] += sz[u];
  mx[x] = max(mx[x], sz[u]);
  tmp.push_back(x);
```

```
void DfsDist(int x, int64_t D = 0) {
 Dist[x].push_back(D); vis[x] = true;
  for (auto [u, w] : g[x])
   if (not vis[u]) DfsDist(u, D + w);
 void DfsCen(int x, int D = 0, int p = -1) {
  tmp.clear(); DfsSz(x);
  int M = tmp.size();
  int C = -1;
  for (int u : tmp) {
  if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
   vis[u] = false;
  DfsDist(C);
  for (int u : tmp) vis[u] = false;
Pa[C] = p; vis[C] = true; Dep[C] = D;
for (auto [u, w] : g[C])
   if (not vis[u]) DfsCen(u, D + 1, C);
 Centroid(int N, vector<vector<pair<int,int>>> g)
  : Sub(N), Sub2(N), Cnt(N), Cnt2(N), Dist(N),
  Pa(N), Dep(N), vis(N), sz(N), mx(N)
  { DfsCen(0); }
 void Mark(int v) {
  int x = v, z = -1;
for (int i = Dep[v]; i >= 0; --i) {
   Sub[x] += Dist[v][i]; Cnt[x]++;
   if (z != -1) {
    Sub2[z] += Dist[v][i];
    Cnt2[z]++;
   z = x; x = Pa[x];
  }
 int64_t Query(int v) {
 int64_t res = 0;
  int x = v, z = -1
  for (int i = Dep[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Cnt[x] * Dist[v][i];
   if (z != -1) res-=Sub2[z]+1LL*Cnt2[z]*Dist[v][i];
   z = x; x = Pa[x];
  return res;
};
      Directed Minimum Spanning Tree
struct Edge { int u, v, w; };
struct DirectedMST { // find maximum
```

```
int solve(vector<Edge> E, int root, int n) {
 int ans = 0;
 while (true) {
  // find best in edge
  vector<int> in(n, -inf), prv(n, -1);
  for (auto e : E)
   if (e.u != e.v && e.w > in[e.v]) {
    in[e.v] = e.w;
    prv[e.v] = e.u;
  in[root] = 0; prv[root] = -1;
  for (int i = 0; i < n; i++)</pre>
  if (in[i] == -inf) return -inf;
  // find cycle
  int tot = 0;
  vector<int> id(n, -1), vis(n, -1);
  for (int i = 0; i < n; i++) {</pre>
   ans += in[i];
   for (int x = i; x != -1 && id[x] == -1; x = prv[x])
    if (vis[x] == i) {
     for (int y = prv[x]; y != x; y = prv[y])
      id[y] = tot;
     id[x] = tot++;
     break;
    vis[x] = i;
  if (!tot) return ans;
  for (int i = 0; i < n; i++)</pre>
   if (id[i] == -1) id[i] = tot++;
```

```
for (auto &e : E) {
   if (id[e.u] != id[e.v]) e.w -= in[e.v];
   e.u = id[e.u], e.v = id[e.v];
}
   n = tot; root = id[root];
}
}
DMST;
```

3.6 Dominator Tree

```
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
 // vertices are numbered from 0 to n-1
 fill(dfn, dfn + n, -1);fill(rev, rev + n, -1);
 fill(fa, fa + n, -1); fill(val, val + n, -1);
 fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
 fill(dom, dom + n, -1); tk = 0;
 for (int i = 0; i < n; ++i) {</pre>
  g[i].clear(); r[i].clear(); rdom[i].clear();
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
 rev[dfn[x] = tk] = x;
 fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
 for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
  r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
 if (fa[x] == x) return c ? -1 : x;
 int p = find(fa[x], 1);
 if (p == -1) return c ? fa[x] : val[x]
 if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
 fa[x] = p;
 return c ? p : val[x];
vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
// p[i] = -2 if i is unreachable from s
 dfs(s);
 for (int i = tk - 1; i >= 0; --i) {
  for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
  if (i) rdom[sdom[i]].push_back(i);
  for (int &u : rdom[i]) {
   int p = find(u);
   if (sdom[p] == i) dom[u] = i;
   else dom[u] = p;
  if (i) merge(i, rp[i]);
 vector<int> p(n, -2); p[s] = -1;
 for (int i = 1; i < tk; ++i)
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];</pre>
 for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
 return p;
}}
```

3.7 Edge Coloring

if (p) X[u] = X[v] = p;

```
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
   for (int i = 0; i <= N; i++)
      for (int j = 0; j <= N; j++)
      C[i][j] = G[i][j] = 0;
}
void solve(vector<pair<int, int>> &E, int N) {
   int X[kN] = {}, a;
   auto update = [&](int u) {
   for (X[u] = 1; C[u][X[u]]; X[u]++);
   };
   auto color = [&](int u, int v, int c) {
      int p = G[u][v];
   G[u][v] = G[v][u] = c;
   C[u][c] = v, C[v][c] = u;
   C[u][p] = C[v][p] = 0;
```

```
else update(u), update(v);
  return p:
 }:
 auto flip = [&](int u, int c1, int c2) {
  int p = C[u][c1];
  swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
if (!C[u][c1]) X[u] = c1;
  if (!C[u][c2]) X[u] = c2;
  return p;
 for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {
    auto [u, v] = E[t];
}</pre>
  int v0 = v, c = X[u], c0 = c, d;
  vector<pair<int, int>> L; int vst[kN] = {};
while (!G[u][v0]) {
   L.emplace_back(v, d = X[v]);
if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
      c = color(u, L[a].first, c);
    else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
      color(u, L[a].first, L[a].second);
    else if (vst[d]) break
   else vst[d] = 1, v = C[u][d];
  if (!G[u][v0]) {
  for (; v; v = flip(v, c, d), swap(c, d));
   if (C[u][c0]) { a = int(L.size()) - 1;
    while (--a >= 0 && L[a].second != c);
     for(;a>=0;a--)color(u,L[a].first,L[a].second);
    } else t--;
}
```

Lowbit Decomposition

```
class LBD {
 int timer, chains;
 vector<vector<int>> G;
 vector<int> tl, tr, chain, head, dep, pa;
 // chains : number of chain
 // tl, tr[u] : subtree interval in the seq. of u
 // head[i] : head of the chain i
 // chian[u] : chain id of the chain u is on
 void predfs(int u, int f) {
  dep[u] = dep[pa[u] = f] + 1;
  for (int v : G[u]) if (v != f) {
   predfs(v, u);
   if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
    chain[u] = chain[v];
  if (chain[u] == 0) chain[u] = ++chains;
 }
 void dfschain(int u, int f) {
  tl[u] = timer++;
  if (head[chain[u]] == -1)
   head[chain[u]] = u;
  for (int v : G[u])
  if (v != f and chain[v] == chain[u])
  dfschain(v, u);
for (int v : G[u])
if (v != f and chain[v] != chain[u])
    dfschain(v, u);
  tr[u] = timer;
 LBD(int n) : timer(0), chains(0), G(n), tl(n), tr(n),
  \begin{array}{c} chain(n)\,,\;head(n,\stackrel{'}{-}1)\,,\;dep(n)\,,\;pa(n)\,\;\{\}\\ void\;add\_edge(int\;u,\;int\;v)\;\;\{ \end{array} 
  G[u].push_back(v); G[v].push_back(u);
 }
 void decompose() { predfs(0, 0); dfschain(0, 0); }
 PII get_subtree(int u) { return {tl[u], tr[u]}; }
 vector<PII> get_path(int u, int v) {
  vector<PII> res;
  while (chain[u] != chain[v]) {
   if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
    swap(u, v)
   int s = head[chain[u]];
   res.emplace_back(tl[s], tl[u] + 1);
   u = pa[s];
```

```
if (dep[u] < dep[v]) swap(u, v);</pre>
  res.emplace_back(tl[v], tl[u] + 1);
  return res;
};
```

3.9 Manhattan Minimum Spanning Tree

```
typedef Point<int> P;
vector<array<int, 3>> manhattanMST(vector<P> ps) {
 vi id(sz(ps));
 iota(all(id), 0);
 vector<array<int, 3>> edges;
 rep(k, 0, 4) {
  sort(all(id),
                [&](int i, int j) {
   return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
  });
  map<int, int> sweep;
for (int i : id) {
   for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++)) {
    int j = it->second;
    P d = ps[i] - ps[j];
if (d.y > d.x) break;
    edges.push_back({d.y + d.x, i, j});
   sweep[-ps[i].y] = i;
  for (P &p : ps)
   if (k \& 1) p.x = -p.x;
   else swap(p.x, p.y);
 return edges; // [{w, i, j}, ...]
```

3.10 MaxClique

```
// contain a self loop u to u, than u won't in clique
template < size_t MAXN >
class MaxClique{
private:
 using bits = bitset< MAXN >;
 bits popped, G[ MAXN ], ans;
 size_t deg[ MAXN ], deo[ MAXN ], n;
 void sort_by_degree() {
  popped.reset();
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
    deg[ i ] = G[ i ].count();
  for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
    size_t mi = MAXN, id = 0;
    for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )</pre>
        mi = deg[id = j];
    popped[ deo[ i ] = id ] = 1;
    for( size_t u = G[ i ]._Find_first() ;
     u < n ; u = G[ i ]._Find_next( u ) )
       -- deg[ u ];
 void BK( bits R, bits P, bits X ) {
  if (R.count()+P.count() <= ans.count()) return;</pre>
  if ( not P.count() and not X.count() )
   if ( R.count() > ans.count() ) ans = R;
   return:
  /* greedily chosse max degree as pivot
  bits cur = P | X; size_t pivot = 0, sz = 0;
  for ( size_t u = cur._Find_first() ;
   u < n ; u = cur._Find_next( u )
    if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
  cur = P & ( ~G[ pivot ] );
  */ // or simply choose first
  bits cur = P & (~G[ ( P | X )._Find_first() ]);
  for ( size_t u = cur._Find_first()
   u < n ; u = cur._Find_next( u ) ) {
   if ( R[ u ] ) continue;
   R[u] = 1
   BK( R, P & G[ u ], X & G[ u ] );
   R[u] = P[u] = 0, X[u] = 1;
public:
 void init( size_t n_ ) {
 n = n_{-};
```

int solve(bitset<kN> mask) { // vertex mask

```
for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                 vector<int> r,
   G[ i ].reset();
                                                                 for (int i = 0; i < n; i++)
 ans.reset();
                                                                 if (mask[i]) r.push_back(i);
                                                                 for (int i = 0; i < n; i++)
                                                                 d[i] = int((a[i] & mask).count());
void add_edges( int u, bits S ) { G[ u ] = S; }
 void add_edge( int u, int v ) {
                                                                 sort(r.begin(), r.end(),
 G[u][v] = G[v][u] = 1;
                                                                  [&](int i, int j) { return d[i] > d[j]; });
                                                                 csort(r, c);
 int solve() {
                                                                 dfs(r, c, 1, mask);
 sort_by_degree(); // or simply iota( deo... )
for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                 return ans; // sol[0 ~ ans-1]
  deg[ i ] = G[ i ].count();
                                                              } graph;
 bits pob, nob = 0; pob.set();
                                                              3.12 Minimum Mean Cycle
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
 for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
                                                              /* minimum mean cycle O(VE) */
   size_t v = deo[ i ];
                                                              struct MMC{
   bits tmp; tmp[v] = 1;
                                                              #define FZ(n) memset((n),0,sizeof(n))
  BK( tmp, pob & G[ v ], nob & G[ v ] );
pob[ v ] = 0, nob[ v ] = 1;
                                                              #define E 101010
                                                              #define V 1021
                                                              #define inf 1e9
  return static_cast< int >( ans.count() );
                                                               struct Edge { int v,u; double c; };
                                                               int n, m, prv[V][V], prve[V][V], vst[V];
};
                                                               Edge e[E];
                                                               vector<int> edgeID, cycle, rho;
3.11
      MaxCliqueDyn
                                                               double d[V][V];
constexpr int kN = 150;
                                                               void init( int _n ) { n = _n; m = 0; }
struct MaxClique { // Maximum Clique
                                                               // WARNING: TYPE matters
bitset<kN> a[kN], cs[kN];
                                                               void add_edge( int vi , int ui , double ci )
int ans, sol[kN], q, cur[kN], d[kN], n;
                                                               { e[ m ++ ] = { vi , ui , ci }; }
void init(int _n) {
                                                               void bellman_ford() {
 n = n, ans q = 0;
                                                                 for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {</pre>
 for (int i = 0; i < n; i++) a[i].reset();</pre>
                                                                  fill(d[i+1], d[i+1]+n, inf);
for(int j=0; j<m; j++) {
void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
                                                                               j<m; j++)
void csort(vector<int> &r, vector<int> &c) {
                                                                   int v = e[j].v, u = e[j].u;
 int mx = 1, km = max(ans - q + 1, 1), t = 0,
                                                                   if(d[i][v]<inf \&\& d[i+1][u]>d[i][v]+e[j].c) {
    m = int(r.size())
                                                                    d[i+1][u] = d[i][v]+e[j].c;
  cs[1].reset(); cs[2].reset();
                                                                    prv[i+1][u] = v;
 for (int i = 0; i < m; i++) {
  int p = r[i], k = 1;</pre>
                                                                    prve[i+1][u] = j;
   while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
   cs[k][p] = 1;
   if (k < km) r[t++] = p;
                                                               double solve(){
                                                                 // returns inf if no cycle, mmc otherwise
  c.resize(m);
                                                                 double mmc=inf;
 if (t) c[t - 1] = 0;
                                                                 int st = -1
 for (int k = km; k <= mx; k++) {</pre>
                                                                 bellman_ford();
  for (int p = int(cs[k]._Find_first());
                                                                 for(int i=0; i<n; i++) {</pre>
      p < kN; p = int(cs[k]._Find_next(p))) {</pre>
                                                                  double avg=-inf;
    r[t] = p; c[t++] = k;
                                                                  for(int k=0; k<n; k++) {</pre>
                                                                   if(d[n][i]<inf-eps)</pre>
  }
                                                                    avg=max(avg,(d[n][i]-d[k][i])/(n-k));
                                                                   else avg=max(avg,inf);
 void dfs(vector<int> &r, vector<int> &c, int 1,
 bitset<kN> mask) {
                                                                  if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
  while (!r.empty()) {
  int p = r.back(); r.pop_back();
                                                                FZ(vst);edgeID.clear();cycle.clear();rho.clear();
   mask[p] = 0;
                                                                 for (int i=n; !vst[st]; st=prv[i--][st]) {
   if (q + c.back() <= ans) return;</pre>
                                                                 vst[st]++;
   cur[q++] = p;
                                                                  edgeID.PB(prve[i][st]);
   vector<int> nr, nc;
                                                                  rho.PB(st);
   bitset<kN> nmask = mask & a[p];
   for (int i : r)
                                                                 while (vst[st] != 2) {
    if (a[p][i]) nr.push_back(i);
                                                                 int v = rho.back(); rho.pop_back();
   if (!nr.empty()) {
                                                                  cycle.PB(v);
    if (1 < 4) {
                                                                 vst[v]++;
     for (int i : nr)
      d[i] = int((a[i] & nmask).count());
                                                                 reverse(ALL(edgeID));
     sort(nr.begin(), nr.end(),
                                                                edgeID.resize(SZ(cycle));
      [&](int x, int y) {
  return d[x] > d[y];
                                                                 return mmc;
      });
                                                              } mmc;
                                                                      Mo's Algorithm on Tree
                                                              3.13
    csort(nr, nc); dfs(nr, nc, l + 1, nmask);
  } else if (q > ans) {
   ans = q; copy(cur, cur + q, sol);
                                                              dfs u:
                                                               push u
   c.pop_back(); q--;
                                                               iterate subtree
 }
                                                               push u
                                                              Let P = LCA(u, v) with St(u) <= St(v)
```

if (P == u) query[St(u), St(v)]

```
else query[Ed(u), St(v)], query[St(P), St(P)]
3.14 Virtual Tree
vector<pair<int, int>> build(vector<int> vs, int r) {
 vector<pair<int, int>> res;
sort(vs.begin(), vs.end(), [](int i, int j) {
  return dfn[i] < dfn[j]; });</pre>
 vector<int> s = {r};
for (int v : vs) if (v != r) {
  if (int o = lca(v, s.back()); o != s.back()) {
   while (s.size() >= 2) {
   if (dfn[s[s.size() - 2]] < dfn[o]) break</pre>
     res.emplace_back(s[s.size() - 2], s.back());
     s.pop_back();
    if (s.back() != o) {
     res.emplace_back(o, s.back());
     s.back() = o;
  s.push_back(v);
 for (size_t i = 1; i < s.size(); ++i)</pre>
  res.emplace_back(s[i - 1], s[i]);
 return res; // (x, y): x->y
      Matching & Flow
4
4.1 Bipartite Matching
struct BipartiteMatching {
```

```
vector<int> X[N];
int fX[N], fY[N], n;
bitset<N> vis
bool dfs(int x)
 for (auto i : X[x]) if (not vis[i]) {
  vis[i] = true;
  if (fY[i] == -1 || dfs(fY[i])) {
   fY[fX[x] = i] = x;
   return true:
 }
 return false;
}
void init(int n_, int m) {
 vis.reset();
fill(X, X + (n = n_), vector<int>());
memset(fX, -1, sizeof(int) * n);
memset(fY, -1, sizeof(int) * m);
void add_edge(int x, int y) { X[x].push_back(y); }
int solve() { // return how many pair matched
 int cnt = 0;
for (int i = 0; i < n; i++) {
 vis.reset();
  cnt += dfs(i);
 return cnt;
```

4.2 Dijkstra Cost Flow

```
// kN = \#(vertices)
// MCMF.{Init, AddEdge, MincostMaxflow}
// MincostMaxflow(source, sink, flow_limit, &cost)
// => flow
using Pii = pair<int, int>;
constexpr int kInf = 0x3f3f3f3f, kN = 500;
struct Edge {
int to, rev, cost, flow;
struct MCMF { // 0-based
int n{}, m{}, s{}, t{};
vector<Edge> graph[kN];
// Larger range for relabeling
int64_t dis[kN] = {}, h[kN] = {};
int p[kN] = {};
void Init(int nn) {
 n = nn;
 for (int i = 0; i < n; i++) graph[i].clear();</pre>
```

```
void AddEdge(int u, int v, int f, int c) {
  graph[u].push_back({v,
   static_cast<int>(graph[v].size()), c, f});
  graph[v].push_back(
   {u, static_cast<int>(graph[u].size()) - 1,
    -c, 0});
 bool Dijkstra(int &max_flow, int64_t &cost) {
  priority_queue<Pii, vector<Pii>, greater<>> pq;
  fill_n(dis, n, kInf);
  dis[s] = 0
  pq.emplace(0, s);
  while (!pq.empty()) {
   auto u = pq.top();
   pq.pop();
   int v = u.second;
   if (dis[v] < u.first) continue;</pre>
   for (auto &e : graph[v]) {
    auto new_dis =
     dis[v] + e.cost + h[v] - h[e.to];
    if (e.flow > 0 && dis[e.to] > new_dis) {
     dis[e.to] = new_dis;
     p[e.to] = e.rev;
     pq.emplace(dis[e.to], e.to);
   }
  if (dis[t] == kInf) return false;
  for (int i = 0; i < n; i++) h[i] += dis[i];</pre>
  int d = max_flow;
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   d = min(d, graph[e.to][e.rev].flow);
  max_flow -= d;
  cost += int64_t(d) * h[t];
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   e.flow += d
   graph[e.to][e.rev].flow -= d;
  }
  return true;
 int MincostMaxflow(
  int ss, int tt, int max_flow, int64_t &cost) {
  this->s = ss, this->t = tt;
  cost = 0;
  fill_n(h, n, 0);
  auto orig_max_flow = max_flow;
  while (Dijkstra(max_flow, cost) && max_flow) {}
  return orig_max_flow - max_flow;
};
4.3 Dinic
```

```
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
    int to, rev;
    Cap cap:
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
      int u = bfs.front(); bfs.pop();
      for (auto e: G[u]) {
        if (e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
        bfs.push(e.to); lv[e.to] = lv[u] + 1;
      }
    return lv[ed] != -1;
  Cap DFS(int u, Cap f){
    if (u == ed) return f;
```

```
Cap ret = 0;
     for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
        auto &e = G[u][i];
        if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
       Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
       if (f == 0) return ret;
     if (ret == 0) lv[u] = -1;
     return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
  G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
  st = st_, ed = ed_; Cap ret = 0;
     while (BFS()) {
        idx.assign(n, 0);
        Cap f = DFS(st, numeric_limits<Cap>::max());
       ret += f;
       if (f == 0) break;
     return ret;
};
```

4.4 Construct VC

```
vi cover(vector<vi>& g, int n, int m) {
  vi match(m, -1);
  int res = dfsMatching(g, match);
  vector<bool> lfound(n, true), seen(m);
  for (int it : match) if (it != -1) lfound[it] = false
 vi q, cover;
  rep(i,0,n) if (lfound[i]) q.push_back(i);
  while (!q.empty()) {
    int i = q.back(); q.pop_back();
    lfound[i] = 1;
    for (int e : g[i]) if (!seen[e] && match[e] != -1)
      seen[e] = true;
      q.push_back(match[e]);
  }
  rep(i,0,n) if (!lfound[i]) cover.push_back(i);
  rep(i,0,m) if (seen[i]) cover.push_back(n+i);
  assert(sz(cover) == res);
  return cover;
```

4.5 Flow Models

- · Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, con- $\mathsf{nect}\ v \to T \ \mathsf{with}\ \mathsf{capacity}\ -in(v).$
 - To maximize, connect $t\, o\,s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Oth-
 - erwise, the maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect $t\, o\,s$ with capacity ∞ and let the flow from S to Tbe f'. If $f + f' \neq \sum_{v \in V, in(v) > 0}^{n} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- · Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M, x \to y$ otherwise. 2. DFS from unmatched vertices in X.

 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source S and sink T2. For each edge (x,y,c), connect $x \to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y \to x$ with (cost,cap)=(-c,1)

- 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
- 4. For each vertex v with d(v)>0, connect S o v with (cost, cap)=0(0, d(v))
- 5. For each vertex v with d(v) < 0, connect $v \rightarrow T$ with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \to v, v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with
 - 5. For $v \in G$, connect it with sink $v \to t$ with capacity K + 2T - $(\sum_{e \in E(v)} w(e)) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with
 - weight w(u,v). 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- · 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x, t) with capacity c_x and create edge (s, y) with
- 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4.6 General Graph Matching

```
namespace matching {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
 for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
 g[u].push_back(v);
 g[v].push_back(u);
int Find(int u) {
 return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
  static int tk = 0; tk++;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
void Blossom(int x, int y, int 1) {
 while (Find(x) != 1) {
  pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
  if (fa[x] == x) fa[x] = 1;
  if (fa[y] == y) fa[y] = 1;
  x = pre[y];
bool Bfs(int r, int n) {
 for (int i = 0; i \le n; ++i) fa[i] = i, s[i] = -1;
 while (!q.empty()) q.pop();
 q.push(r);
 s[r] = 0;
 while (!q.empty()) {
  int x = q.front(); q.pop();
  for (int u : g[x]) {
```

}

```
if (s[u] == -1) {
                                                                return rt;
    pre[u] = x, s[u] = 1;
    if (match[u] == n) {
                                                               4.9 Kuhn Munkres
     for (int a = u, b = x, last; b != n; a = last, b =
     pre[a])
                                                               class KM {
      last = match[b], match[b] = a, match[a] = b;
                                                               private:
     return true;
                                                                static constexpr lld INF = 1LL << 60;</pre>
                                                                vector<lld> hl,hr,slk;
    q.push(match[u]);
                                                                vector<int> fl,fr,pre,qu;
    s[match[u]] = 0;
                                                                vector<vector<lld>> w;
   } else if (!s[u] && Find(u) != Find(x)) {
                                                                vector<bool> v1.vr:
    int 1 = LCA(u, x, n);
Blossom(x, u, 1);
Blossom(u, x, 1);
                                                                int n, ql, qr;
                                                                bool check(int x) {
                                                                 if (v1[x] = true, f1[x] != -1)
                                                                  return vr[qu[qr++] = f1[x]] = true;
                                                                 while (x != -1) swap(x, fr[fl[x] = pre[x]]);
                                                                 return false:
 return false;
                                                                void bfs(int s) {
int Solve(int n) {
                                                                 fill(slk.begin(), slk.end(), INF);
                                                                 fill(v1.begin(), v1.end(), false);
 int res = 0;
 for (int x = 0; x < n; ++x) {
                                                                 fill(vr.begin(), vr.end(), false);
  if (match[x] == n) res += Bfs(x, n);
                                                                 ql = qr = 0;
                                                                 vr[qu[qr++] = s] = true;
 return res;
                                                                 while (true) {
                                                                  11d d;
                                                                  while (ql < qr) {</pre>
4.7 Global Min-Cut
                                                                   for (int x = 0, y = qu[ql++]; x < n; ++x) {
const int maxn = 500 + 5
                                                                    if(!vl[x]&&slk[x]>=(d=hl[x]+hr[y]-w[x][y])){
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
                                                                     if (pre[x] = y, d) slk[x] = d;
                                                                     else if (!check(x)) return;
void add_edge(int x, int y, int c) {
                                                                    }
w[x][y] += c; w[y][x] += c;
                                                                   }
                                                                  d = INF;
pair<int, int> phase(int n) {
 memset(v, false, sizeof(v));
                                                                  for (int x = 0; x < n; ++x)
 memset(g, 0, sizeof(g));
                                                                   if (!v1[x] && d > slk[x]) d = slk[x];
 int s = -1, t = -1;
                                                                  for (int x = 0; x < n; ++x) {
 while (true) {
                                                                   if (v1[x]) h1[x] += d;
                                                                   else slk[x] -= d;
  int c = -1:
  for (int i = 0; i < n; ++i) {</pre>
                                                                   if (vr[x]) hr[x] -= d;
   if (del[i] || v[i]) continue;
   if (c == -1 \mid | g[i] > g[c]) c = i;
                                                                  for (int x = 0; x < n; ++x)
                                                                   if (!v1[x] && !slk[x] && !check(x)) return;
  if (c == -1) break;
                                                                 }
  v[s = t, t = c] = true;
  for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
                                                               public:
                                                                void init( int n_ ) {
   g[i] += w[c][i];
                                                                 qu.resize(n = n_);
                                                                 fl.assign(n, -1); fr.assign(n, -1);
  }
                                                                 hr.assign(n, 0); hl.resize(n);
 return make_pair(s, t);
                                                                 w.assign(n, vector<lld>(n));
                                                                 slk.resize(n); pre.resize(n);
int mincut(int n) {
                                                                 vl.resize(n); vr.resize(n);
 int cut = 1e9;
 memset(del, false, sizeof(del));
                                                                void set_edge( int u, int v, lld x ) {w[u][v] = x;}
                                                                11d solve() {
 for (int i = 0; i < n - 1; ++i) {
 int s, t; tie(s, t) = phase(n);
del[t] = true; cut = min(cut, g[t]);
                                                                 for (int i = 0; i < n; ++i)</pre>
                                                                  hl[i] = *max_element(w[i].begin(), w[i].end());
  for (int j = 0; j < n; ++j) {
                                                                 for (int i = 0; i < n; ++i) bfs(i);</pre>
                                                                 11d res = 0;
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
                                                                 for (int i = 0; i < n; ++i) res += w[i][f1[i]];</pre>
  }
                                                                 return res;
 return cut;
                                                                }
                                                               } km;
                                                                       Minimum Cost Circulation
4.8 GomoryHu Tree
                                                               struct Edge { int to, cap, rev, cost; };
int g[maxn];
vector<edge> GomoryHu(int n){
                                                               vector<Edge> g[kN];
                                                               int dist[kN], pv[kN], ed[kN];
 vector<edge> rt;
 for(int i=1;i<=n;++i)g[i]=1;</pre>
                                                               bool mark[kN];
 for(int i=2;i<=n;++i){</pre>
                                                               int NegativeCycle(int n) {
                                                                memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
  int t=g[i];
  flow.reset(); // clear flows on all edge
rt.push_back({i,t,flow(i,t)});
                                                                int upd = -1;
                                                                for (int i = 0; i <= n; ++i) {</pre>
  flow.walk(i); // bfs points that connected to i (use
                                                                 for (int j = 0; j < n; ++j) {
    edges not fully flow)
  for(int j=i+1;j<=n;++j){</pre>
                                                                  int idx = 0;
  if(g[j]==t && flow.connect(j))g[j]=i; // check if i
                                                                  for (auto &e : g[j]) {
                                                                   if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
    can reach i
                                                                    dist[e.to] = dist[j] + e.cost;
```

pv[e.to] = j, ed[e.to] = idx;

```
if (i == n) {
      upd = j;
      while(!mark[upd])mark[upd]=1,upd=pv[upd];
      return upd;
    idx++;
  }
 return -1:
int Solve(int n) {
int rt = -1, ans = 0;
while ((rt = NegativeCycle(n)) >= 0) {
 memset(mark, false, sizeof(mark));
 vector<pair<int, int>> cyc;
 while (!mark[rt]) {
  cyc.emplace_back(pv[rt], ed[rt]);
  mark[rt] = true;
  rt = pv[rt];
  }
  reverse(cyc.begin(), cyc.end());
  int cap = kInf;
  for (auto &i : cyc)
  auto &e = g[i.first][i.second];
   cap = min(cap, e.cap);
 for (auto &i : cyc) {
  auto &e = g[i.first][i.second];
   e.cap -= cap;
   g[e.to][e.rev].cap += cap;
   ans += e.cost * cap;
return ans;
```

4.11 Minimum Cost Maximum Flow

```
class MiniCostMaxiFlow{
using Cap = int; using Wei = int64_t;
using PCW = pair<Cap,Wei>;
static constexpr Cap INF_CAP = 1 << 30;</pre>
static constexpr Wei INF_WEI = 1LL<<60;</pre>
private:
struct Edge{
 int to, back;
 Cap cap; Wei wei;
 Edge() {}
 Edge(int a,int b, Cap c, Wei d):
  to(a),back(b),cap(c),wei(d) {}
};
int ori, edd;
vector<vector<Edge>> G;
vector<int> fa, wh;
vector<bool> inq;
vector<Wei> dis;
PCW SPFA(){
 fill(inq.begin(),inq.end(),false);
 fill(dis.begin(), dis.end(), INF_WEI);
 queue<int> qq; qq.push(ori);
 dis[ori] = 0;
 while(not qq.empty()){
  int u=qq.front();qq.pop();
   inq[u] = false;
   for(int i=0;i<SZ(G[u]);++i){</pre>
   Edge e=G[u][i];
    int v=e.to; Wei d=e.wei;
    if(e.cap<=0||dis[v]<=dis[u]+d)</pre>
     continue:
    dis[v] = dis[u] + d;
    fa[v] = u, wh[v] = i;
    if (inq[v]) continue;
    qq.push(v);
    inq[v] = true;
 if(dis[edd]==INF_WEI) return {-1, -1};
 Cap mw=INF_CAP;
 for(int i=edd;i!=ori;i=fa[i])
  mw=min(mw,G[fa[i]][wh[i]].cap);
 for (int i=edd;i!=ori;i=fa[i]){
```

```
eg.cap -= mw;
   G[eg.to][eg.back].cap+=mw;
  return {mw, dis[edd]};
public:
 void init(int n){
  G.clear();G.resize(n);
  fa.resize(n);wh.resize(n);
  inq.resize(n); dis.resize(n);
 void add_edge(int st, int ed, Cap c, Wei w){
  G[st].emplace_back(ed,SZ(G[ed]),c,w);
  G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
 PCW solve(int a, int b){
  ori = a, edd = b;
  Cap cc=0; Wei ww=0;
  while(true){
   PCW ret=SPFA();
   if(ret.first==-1) break;
   cc+=ret.first;
   ww+=ret.first * ret.second:
  return {cc.ww}:
 }
} mcmf;
```

4.12 Maximum Weight Graph Matching

```
struct WeightGraph {
 static const int inf = INT_MAX;
 static const int maxn = 514;
 struct edge {
  int u, v, w;
  edge(){}
  edge(int u, int v, int w): u(u), v(v), w(w) {}
 int n, n_x;
 edge g[maxn * 2][maxn * 2];
 int lab[maxn * 2];
 int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa
    [maxn * 2]
 int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
    maxn * 2];
 vector<int> flo[maxn * 2];
 queue<int> q;
 int e_delta(const edge &e) { return lab[e.u] + lab[e.v
    ] - g[e.u][e.v].w * 2; }
 void update_slack(int u, int x) { if (!slack[x] ||
    e_delta(g[u][x]) < e_delta(g[slack[x]][x])) slack[x
    ] = u; }
 void set_slack(int x) {
  slack[x] = 0;
  for (int u = 1; u <= n; ++u)</pre>
   if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
    update_slack(u, x);
 void q_push(int x) -
 if (x \le n) q.push(x);
  else for (size_t i = 0; i < flo[x].size(); i++)</pre>
    q_push(flo[x][i]);
 void set_st(int x, int b) {
  st[x] = b;
  if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)</pre>
     set_st(flo[x][i], b);
 int get_pr(int b, int xr) {
  int pr = find(flo[b].begin(), flo[b].end(), xr) - flo
    [b].begin();
  if (pr % 2 == 1) {
   reverse(flo[b].begin() + 1, flo[b].end());
   return (int)flo[b].size() - pr;
  return pr;
 void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u, xr);
```

```
for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo</pre>
   [u][i ^ 1]);
 set_match(xr, v);
 rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
   end());
void augment(int u, int v) {
for (; ; ) {
  int xnv = st[match[u]];
  set_match(u, v);
  if (!xnv) return;
  set_match(xnv, st[pa[xnv]]);
  u = st[pa[xnv]], v = xnv;
 }
int get_lca(int u, int v) {
 static int t = 0;
 for (++t; u || v; swap(u, v)) {
  if (u == 0) continue;
  if (vis[u] == t) return u;
  vis[u] = t;
  u = st[match[u]];
  if (u) u = st[pa[u]];
 return 0;
void add_blossom(int u, int lca, int v) {
 int b = n + 1;
 while (b <= n_x && st[b]) ++b;</pre>
 if (b > n_x) ++n_x;
lab[b] = 0, S[b] = 0;
 match[b] = match[lca];
 flo[b].clear();
 flo[b].push_back(lca);
 for (int x = u, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
 reverse(flo[b].begin() + 1, flo[b].end())
 for (int x = v, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
 set_st(b, b);
 for (int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].w
   = 0;
 for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  for (int x = 1; x <= n_x; ++x)
   if (g[b][x].w == 0 \mid \mid e_{delta}(g[xs][x]) < e_{delta}(g[xs][x])
   [b][x]))
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
  for (int x = 1; x <= n; ++x)
   if (flo_from[xs][x]) flo_from[b][x] = xs;
 set_slack(b);
void expand_blossom(int b) {
for (size_t i = 0; i < flo[b].size(); ++i)
set_st(flo[b][i], flo[b][i]);</pre>
 int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
   xr):
 for (int i = 0; i < pr; i += 2)</pre>
  int xs = flo[b][i], xns = flo[b][i + 1];
  pa[xs] = g[xns][xs].u;
  S[xs] = 1, S[xns] = 0;
  slack[xs] = 0, set_slack(xns);
  q_push(xns);
 S[xr] = 1, pa[xr] = pa[b];
 for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  S[xs] = -1, set_slack(xs);
 st[b] = 0;
bool on_found_edge(const edge &e) {
 int u = st[e.u], v = st[e.v];
 if (S[v] == -1) {
  pa[v] = e.u, S[v] = 1;
  int nu = st[match[v]];
slack[v] = slack[nu] = 0;
  S[nu] = 0, q_push(nu);
```

```
} else if (S[v] == 0) {
  int lca = get_lca(u, v);
  if (!lca) return augment(u,v), augment(v,u), true;
  else add_blossom(u, lca, v);
 return false;
bool matching() {
memset(S + 1, -1, sizeof(int) * n_x);
memset(slack + 1, 0, sizeof(int) * n_x);
 q = queue<int>();
 for (int x = 1; x <= n_x; ++x)
 if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0,
   q_push(x);
 if (q.empty()) return false;
 for (; ; )
  while (q.size()) {
   int u = q.front(); q.pop();
   if (S[st[u]] == 1) continue;
   for (int v = 1; v <= n; ++v)
    if (g[u][v].w > 0 && st[u] != st[v]) {
     if (e_delta(g[u][v]) == 0) {
      if (on_found_edge(g[u][v])) return true;
     } else update_slack(u, st[v]);
  int d = inf;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b && S[b] == 1) d = min(d, lab[b] / 2)
  for (int x = 1; x <= n_x; ++x)
   if (st[x] == x && slack[x]) {
    if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x])
    else if (S[x] == 0) d = min(d, e_delta(g[slack[x
   ]][x]) / 2);
  for (int u = 1; u <= n; ++u) {
  if (S[st[u]] == 0) {</pre>
    if (lab[u] <= d) return 0;</pre>
    lab[u] -= d;
   } else if (S[st[u]] == 1) lab[u] += d;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b) {
    if (S[st[b]] == 0) lab[b] += d * 2;
    else if (S[st[b]] == 1) lab[b] -= d * 2;
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
   if (st[x] == x && slack[x] && st[slack[x]] != x &&
   e_delta(g[slack[x]][x]) == 0)
    if (on_found_edge(g[slack[x]][x])) return true;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b && S[b] == 1 && lab[b] == 0)
   expand_blossom(b);
 return false;
pair<long long, int> solve() {
memset(match + 1, 0, sizeof(int) * n);
 n_x = n
 int n_matches = 0;
 long long tot_weight = 0;
 for (int u = 0; u \le n; ++u) st[u] = u, flo[u].clear
   ();
 int w_max = 0;
 for (int u = 1; u <= n; ++u)
  for (int v = 1; v <= n; ++v) {
   flo_from[u][v] = (u == v ? u : 0);
   w_{max} = max(w_{max}, g[u][v].w);
 for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
while (matching()) ++n_matches;
for (int u = 1; u <= n; ++u)</pre>
  if (match[u] && match[u] < u)</pre>
   tot_weight += g[u][match[u]].w;
 return make_pair(tot_weight, n_matches);
void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
   g[vi][ui].w = wi; }
void init(int _n) {
```

```
n = _n;
for (int u = 1; u <= n; ++u)
  for (int v = 1; v <= n; ++v)
    g[u][v] = edge(u, v, 0);
}
};</pre>
```

5 Math

5.1 Common Bounds

5.1.1 Partition function

$$\begin{split} p(0) = 1, \; p(n) = \sum_{k \in \mathbb{Z} \backslash \{0\}} (-1)^{k+1} p(n-k(3k-1)/2) \\ p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n}) \\ \\ \frac{n \quad \mid \; 01234567892050100}{p(n) \quad \mid \; 11235711152230627 \sim & 2e5 \sim & 2e8} \end{split}$$

5.1.2 Divisor function

n	100	1e3	1e6	1e9	1e12	1e15	1e18
$\max_{i < n} (d(i))$	12	32	240	1344	6720	26880	103680

5.1.3 Factorial

n	123	4 !	5 6	7	8		9	10	
n!	126	24 12	20 720	504	0 403	20 362	880 3	62880	00
n						15			
n!	4.0e	7 4.8	e8 6.2	e9 8.	7e10 1	.3e12 2	l.1e13 3	3.6e14	
n	20	25	30	40	50	100	150		171
n!	2e18	2e25	3e32	8e47	' 3e64	9e157	6e26	2 > DE	BL_MAX

5.1.4 Binom Coef

5.2 Strling Number

5.2.1 First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1) \dots (x+n-1) = \sum_{k=0}^{n} S_1(n,k) x^k$$

$$g(x) = x(x+1) \dots (x+n-1) = \sum_{k=0}^{n} a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^{n} \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^{k} ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.2.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k \binom{k}{i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b==(x/y)*a;
}
```

5.4 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
 vector<T> d(output.size() + 1), me, he;
 for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
 for (size_t j = 0; j < me.size(); ++j)
d[i] += output[i - j - 2] * me[j];</pre>
  if ((d[i] -= output[i - 1]) == 0) continue;
  if (me.empty()) {
   me.resize(f = i);
   continue;
  }
  vector<T> o(i - f - 1);
  T k = -d[i] / d[f]; o.push_back(-k);
  for (T x : he) o.push_back(x * k);
  if (o.size() < me.size()) o.resize(me.size());</pre>
  for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
  if (i-f+he.size() >= me.size()) he = me, f = i;
  me = o;
 return me;
```

5.5 Charateristic Polynomial

```
vector<vector<int>> Hessenberg(const vector<vector<int
    >> &A) {
 int N = A.size();
 vector<vector<int>> H = A;
 for (int i = 0; i < N - 2; ++i) {
  if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {
    if (H[j][i]) {
     for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
     for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k
    ][j]);
     break:
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
   * H[i + 1][k] * (kP - coef)) % kP;
for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
 return H;
vector<int> CharacteristicPoly(const vector<vector<int
    >> &A) {
 int N = A.size();
 auto H = Hessenberg(A);
 for (int i = 0; i < N; ++i) {
  for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
 vector<vector<int>> P(N + 1, vector<int>(N + 1));
 P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
  P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
    1];
  int val = 1;
  for (int j = i - 1; j >= 0; --j) {
  int coef = 1LL * val * H[j][i - 1] % kP;
   for (int k = 0; k <= j; ++k) P[i][k] = (P[i][k] + 1
LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 if (N & 1) {
 for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
 return P[N];
```

5.6 Chinese Remainder

```
x = a1 % m1
x = a2 % m2
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
// 0 <= x < lcm(m1, m2)</pre>
```

5.7 De-Bruijn

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
  if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
} else {
 aux[t] = aux[t - p];
  db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {</pre>
  aux[t] = i;
   db(t + 1, t, n, k);
}
int de_bruijn(int k, int n) {
  // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
if (k == 1) {
 res[0] = 0;
  return 1;
for (int i = 0; i < k * n; i++) aux[i] = 0;</pre>
db(1, 1, n, k);
 return sz;
```

5.8 DiscreteLog

```
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
  // x^? \equiv y (mod M)
  Int t = 1, c = 0, g = 1;
  for (Int M_ = M; M_ > 0; M_ >>= 1)
   g = g * x % M;
  for (g = gcd(g, M); t % g != 0; ++c) {
   if (t == y) return c;
t = t * x % M;
  if (y % g != 0) return -1;
 t /= g, y /= g, M /= g;
 Int h = 0, gs = 1;

for (; h * h < M; ++h) gs = gs * x % M;
 unordered_map<Int, Int> bs;
 for (Int s = 0; s < h; bs[y] = ++s)
   y = y * x % M;
  for (Int s = 0; s < M; s += h) {
   t = t * gs % M;
    if (bs.count(t)) return c + s + h - bs[t];
  return -1;
```

5.9 Extended Euler

```
a^b \equiv \begin{cases} a^{(b \mod \varphi(m)) + \varphi(m)} & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m
```

5.10 ExtendedFloorSum

```
\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}
```

```
\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}
```

```
5.11 Fast Fourier Transform
const int mod = 1000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
  static_assert (M1 <= M2 && M2 <= M3);
  constexpr int64_t r12 = modpow(M1, M2-2, M2);
  constexpr int64_t r13 = modpow(M1, M3-2, M3);
  constexpr int64_t r23 = modpow(M2, M3-2, M3);
  constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
  B = (B - A + M2) * r12 % M2;
  C = (C - A + M3) * r13 % M3;
  C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
 for (int i = 0; i <= maxn; i++)</pre>
  omega[i] = cplx(cos(2 * pi * j / maxn),
     sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
 int z = __builtin_ctz(n) - 1;
for (int i = 0; i < n; ++i) {</pre>
  int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^{=(i >> j & 1) << (z - j);
  if (x > i) swap(v[x], v[i]);
 for (int s = 2; s <= n; s <<= 1) {
  int z = s >> 1;
for (int i = 0; i < n; i += s) {</pre>
   for (int k = 0; k < z; ++k) {
    cplx x = v[i + z + k] * omega[maxn / s * k];
    v[i + z + k] = v[i + k] - x;
    v[i+k] = v[i+k] + x;
  }
void ifft(vector<cplx> &v, int n) {
 fft(v, n); reverse(v.begin() + 1, v.end());
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
VL convolution(const VI &a, const VI &b) {
 // Should be able to handle N <= 10^5, C <= 10^4
 int sz = 1;
 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 vector<cplx> v(sz);
 for (int i = 0; i < sz; ++i) {
  double re = i < a.size() ? a[i] : 0;
double im = i < b.size() ? b[i] : 0;</pre>
  v[i] = cplx(re, im);
 fft(v, sz);
 for (int i = 0; i <= sz / 2; ++i) {
  int j = (sz - i) & (sz - 1);
  cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
    * cplx(0, -0.25);
  if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
    ].conj()) * cplx(0, -0.25);
  v[i] = x;
```

ifft(v, sz);

VL c(sz);

```
for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
                                                                            x1), inv = (x0-x1, x1) w/o final div */
                                                              * x = (x0+x1,
                                                            void fwt(int x[], int N, bool inv = false) {
return c:
                                                               for (int d = 1; d < N; d <<= 1) {
VI convolution_mod(const VI &a, const VI &b, int p) {
                                                                 for (int s = 0, d2 = d * 2; s < N; s += d2)
                                                                   for (int i = s, j = s + d; i < s + d; i++, j++) {
int sz = 1:
while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
                                                                     int ta = x[i], tb = x[j];
 vector<cplx> fa(sz), fb(sz);
                                                                     x[i] = modadd(ta, tb);
for (int i = 0; i < (int)a.size(); ++i)</pre>
                                                                     x[j] = modsub(ta, tb);
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 for (int i = 0; i < (int)b.size(); ++i)</pre>
 fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
                                                               if (inv) for (int i = 0, invn = modinv(N); i < N; i</pre>
 fft(fa, sz), fft(fb, sz);
double r = 0.25 / sz;
                                                                 x[i] = modmul(x[i], invn);
 cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
 int j = (sz - i) & (sz - 1);
cplx a1 = (fa[i] + fa[j].conj());
                                                            5.14 Miller Rabin
                                                            bool isprime(llu x) {
  static auto witn = [](llu a, llu u, llu n, int t) {
 cplx a2 = (fa[i] - fa[j].conj()) * r2;
 cplx b1 = (fb[i] + fb[j].conj()) * r3;
                                                               if (!a) return false;
  cplx b2 = (fb[i] - fb[j].conj()) * r4;
                                                               while (t--) {
  if (i != j) {
                                                               llu a2 = mmul(a, a, n);
  cplx c1 = (fa[j] + fa[i].conj());
                                                                if (a2 == 1 && a != 1 && a != n - 1) return true;
   cplx c2 = (fa[j] - fa[i].conj()) * r2;
                                                                a = a2;
   cplx d1 = (fb[j] + fb[i].conj()) * r3;
                                                               }
   cplx d2 = (fb[j] - fb[i].conj()) * r4;
                                                               return a != 1;
   fa[i] = c1 * d1 + c2 * d2 * r5;
   fb[i] = c1 * d2 + c2 * d1;
                                                              if (x < 2) return false;</pre>
                                                              if (!(x & 1)) return x == 2;
  fa[j] = a1 * b1 + a2 * b2 * r5;
                                                              int t = __builtin_ctzll(x - 1);
  fb[j] = a1 * b2 + a2 * b1;
                                                              llu \ odd = (x - 1) >> t;
                                                             for (llu m:
  {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
fft(fa, sz), fft(fb, sz);
vector<int> res(sz);
                                                               if (witn(mpow(m % x, odd, x), odd, x, t))
for (int i = 0; i < sz; ++i) {</pre>
                                                               return false:
 long long a = round(fa[i].re), b = round(fb[i].re),
                                                              return true;
       c = round(fa[i].im);
 res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
}
                                                            5.15 NTT
 return res;
                                                            template <int mod, int G, int maxn>
}}
                                                             struct NTT {
5.12 FloorSum
                                                              static_assert (maxn == (maxn & -maxn));
                                                              int roots[maxn];
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
                                                              NTT () {
                                                               int r = modpow(G, (mod - 1) / maxn);
// @return sum_\{i=\theta\}^{n-1} floor((ai + b)/m) mod 2^64
                                                               for (int i = maxn >> 1; i; i >>= 1) {
1lu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
                                                                roots[i] = 1;
11u ans = 0;
                                                                for (int j = 1; j < i; j++)</pre>
while (true) {
                                                                 roots[i + j] = modmul(roots[i + j - 1], r);
 if (a >= m) {
                                                                r = modmul(r, r);
  ans += n * (n - 1) / 2 * (a / m); a %= m;
  if (b >= m) {
                                                              // n must be 2^k, and 0 \ll F[i] \ll mod
  ans += n * (b / m); b %= m;
                                                              void operator()(int F[], int n, bool inv = false) {
                                                               for (int i = 0, j = 0; i < n; i++) {
 llu y_max = a * n + b;
                                                               if (i < j) swap(F[i], F[j]);</pre>
 if (y_max < m) break;</pre>
                                                                for (int k = n>1; (j^k < k; k>=1);
 // y_max < m * (n + 1)
 // floor(y_max / m) <= n
                                                               for (int s = 1; s < n; s *= 2) {
 n = (11u)(y_max / m), b = (11u)(y_max % m);
                                                                for (int i = 0; i < n; i += s * 2) {
 swap(m, a);
                                                                 for (int j = 0; j < s; j++) {
                                                                  int a = F[i+j]
return ans;
                                                                  int b = modmul(F[i+j+s], roots[s+j]);
                                                                  F[i+j] = modadd(a, b); // a + b
11d floor_sum(lld n, lld m, lld a, lld b) {
                                                                  F[i+j+s] = modsub(a, b); // a - b
llu ans = 0;
if (a < 0) {
 11u a2 = (a \% m + m) \% m;
 ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
                                                               if (inv) {
 a = a2:
                                                               int invn = modinv(n);
                                                                for (int i = 0; i < n; i++)</pre>
if (b < 0) {
                                                                F[i] = modmul(F[i], invn);
 11u b2 = (b \% m + m) \% m;
                                                                reverse(F + 1, F + n);
 ans -= 1ULL * n * ((b2 - b) / m);
 b = b2;
return ans + floor_sum_unsigned(n, m, a, b);
                                                            NTT<2013265921, 31, 1048576> ntt;
5.13 FWT
                                                            5.16 Partition Number
/* or convolution:
                                                            int b = sqrt(n)
* x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                            ans[0] = tmp[0] = 1;
* and convolution:
                                                            for (int i = 1; i <= b; i++) {
```

#define fi(1, r) for (int i = int(1); i < int(r); ++i)

```
for (int rep = 0; rep < 2; rep++)</pre>
                                                              template <int mod, int G, int maxn> struct Poly : V {
  for (int j = i; j <= n - i * i; j++)
                                                               static uint32_t n2k(uint32_t n) {
   modadd(tmp[j], tmp[j-i]);
                                                                if (n <= 1) return 1;
 for (int j = i * i; j <= n; j++)
modadd(ans[j], tmp[j - i * i]);</pre>
                                                                return 1u << (32 - __builtin_clz(n - 1));</pre>
                                                               static NTT<mod, G, maxn> ntt; // coefficients in [0, P)
                                                               explicit Poly(int n = 1) : V(n) {}
      Pi Count (Linear Sieve)
                                                               Poly(const V &v) : V(v) {}
                                                               Poly(const Poly &p, size_t n) : V(n) {
static constexpr int N = 1000000 + 5;
                                                                copy_n(p.data(), min(p.size(), n), data());
11d pi[N];
vector<int> primes;
                                                               Poly &irev() { return reverse(data(), data() + size())
bool sieved[N];
                                                                    *this; }
                                                               Poly &isz(int sz) { return resize(sz), *this; }
11d cube_root(11d x){
 lld s=cbrt(x-static_cast<long double>(0.1));
                                                               Poly &iadd(const Poly &rhs) { // n() == rhs.n()
 while(s*s*s <= x) ++s;</pre>
                                                                fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
 return s-1;
                                                                return *this;
11d square_root(11d x){
                                                               Poly &imul(int k) {
 1ld s=sqrt(x-static_cast<long double>(0.1));
                                                                fi(0, size())(*this)[i] = modmul((*this)[i], k);
 while(s*s <= x) ++s;</pre>
                                                                return *this:
 return s-1;
                                                               Poly Mul(const Poly &rhs) const {
void init(){
                                                                const int sz = n2k(size() + rhs.size() - 1);
primes.reserve(N)
                                                                Poly X(*this, sz), Y(rhs, sz);
ntt(X.data(), sz), ntt(Y.data()
 primes.push_back(1);
                                                                                                   sz):
 for(int i=2;i<N;i++) {</pre>
                                                                fi(0, sz) X[i] = modmul(X[i], Y[i]);
  if(!sieved[i]) primes.push_back(i);
                                                                ntt(X.data(), sz, true);
  pi[i] = !sieved[i] + pi[i-1];
                                                                return X.isz(size() + rhs.size() - 1);
  for(int p: primes) if(p > 1) {
  if(p * i >= N) break;
                                                               Poly Inv() const { // coef[0] != 0
   sieved[p * i] = true;
                                                                if (size() == 1) return V{modinv(*begin())};
   if(p % i == 0) break;
                                                                const int sz = n2k(size() * 2);
                                                                Poly X = Poly(*this, (size() + 1) / 2).Inv().isz(sz),
 }
                                                                    Y(*this, sz);
                                                                ntt(X.data(), sz), ntt(Y.data(), sz);
11d phi(11d m, 11d n) {
 static constexpr int MM = 80000, NN = 500;
                                                                fi(0, sz) X[i] = modmul(X[i], modsub(2, modmul(X[i],
                                                                  Y[i])));
 static lld val[MM][NN];
                                                                ntt(X.data(), sz, true);
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
                                                                return X.isz(size());
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
                                                               Poly Sqrt() const { // coef[0] \in [1, mod)^2
 lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
                                                                if (size() == 1) return V{QuadraticResidue((*this)
 if(m<MM&&n<NN) val[m][n] = ret+1;
                                                                  [0], mod)};
 return ret;
                                                                Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
                                                                  size());
11d pi_count(11d);
                                                                return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
11d P2(11d m, 11d n) {
                                                                  + 1):
 11d sm = square_root(m), ret = 0;
 for(lld i = n+1;primes[i]<=sm;i++)</pre>
                                                               pair<Poly, Poly> DivMod(const Poly &rhs) const {
  ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
                                                                if (size() < rhs.size()) return {V{0}, *this};</pre>
 return ret;
                                                                const int sz = size() - rhs.size() + 1;
                                                                Poly X(rhs); X.irev().isz(sz);
11d pi_count(11d m) {
                                                                Poly Y(*this); Y.irev().isz(sz);
 if(m < N) return pi[m];</pre>
                                                                Poly Q = Y.Mul(X.Inv()).isz(sz).irev();
 11d n = pi_count(cube_root(m));
                                                                X = rhs.Mul(Q), Y = *this;
 return phi(m, n) + n - 1 - P2(m, n);
                                                                fi(0, size()) Y[i] = modsub(Y[i], X[i]);
                                                                return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
       Pollard Rho
                                                               Poly Dx() const {
// does not work when n is prime
                                                                Poly ret(size() - 1);
// return any non-trivial factor
                                                                fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
llu pollard_rho(llu n) {
 static auto f = [](llu x, llu k, llu m) {
                                                                return ret.isz(max<int>(1, ret.size()));
    return add(k, mul(x, x, m), m); };
 if (!(n & 1)) return 2;
                                                               Poly Sx() const {
 mt19937 rnd(120821011);
                                                                Poly ret(size() + 1);
 while (true) {
                                                                fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
 llu y = 2, yy = y, x = rnd() % n, t = 1;
for (llu sz = 2; t == 1; sz <<= 1, y = yy) {
                                                                  this)[i]);
                                                                return ret;
   for (llu i = 0; t == 1 && i < sz; ++i) {
    yy = f(yy, x, n);
                                                               Poly Ln() const { // coef[0] == 1
    t = gcd(yy > y ? yy - y : y - yy, n);
                                                                return Dx().Mul(Inv()).Sx().isz(size());
   }
                                                               Poly Exp() const \{ // coef[\theta] == \theta \}
  if (t != 1 && t != n) return t;
                                                                if (size() == 1) return V{1};
                                                                Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size)
                                                                  ()):
                                                                Poly Y = X.Ln(); Y[0] = mod - 1;
      Polynomial Operations
5.19
                                                                fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
return X.Mul(Y).isz(size());
using V = vector<int>;
```

// maximize c^Tx under Ax <= B

```
Poly Pow(const string &K) const {
                                                                  // return VD(n, -inf) if the solution doesn't exist
                                                                  // return VD(n, +inf) if the solution is unbounded
  int nz = 0;
  while (nz < size() && !(*this)[nz]) ++nz;</pre>
                                                                  using VD = vector<double>;
  int nk = 0, nk2 = 0;
                                                                  using VVD = vector<vector<double>>;
  for (char c : K) {
                                                                  const double eps = 1e-9;
   nk = (nk * 10 + c - '0') % mod;
                                                                  const double inf = 1e+9;
   nk2 = nk2 * 10 + c - '0';
                                                                   int n, m;
   if (nk2 * nz >= size())
                                                                  VVD d:
                                                                  vector<int> p, q;
void pivot(int r, int s) {
    return Poly(size());
   nk2 %= mod - 1;
                                                                    double inv = 1.0 / d[r][s];
  if (!nk && !nk2) return Poly(V{1}, size());
                                                                    for (int i = 0; i < m + 2; ++i)
                                                                    for (int j = 0; j < n + 2; ++j)
if (i != r && j != s)
  Poly X = V(data() + nz, data() + size() - nz * (nk2 - nz)
     1));
  int x0 = X[0];
                                                                       d[i][j] -= d[r][j] * d[i][s] * inv;
                                                                    for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
  return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
    modpow(x0, nk2)).irev().isz(size()).irev();
                                                                    d[r][s] = inv; swap(p[r], q[s]);
 Poly InvMod(int L) { // (to evaluate linear recursion)
  Poly R{1, 0}; // *this * R mod x^L = 1 (*this[0] ==
                                                                  bool phase(int z) {
                                                                   int x = m + z
  for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                                    while (true) {
   Poly 0 = R.Mul(Poly(data(), min<int>(2 << level,
                                                                     int s = -1;
    size())))
                                                                     for (int i = 0; i <= n; ++i) {</pre>
   Poly Q(2 << level); Q[0] = 1;
                                                                      if (!z && q[i] == -1) continue;
   for (int j = (1 << level); j < (2 << level); ++j)</pre>
                                                                      if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
    Q[j] = modsub(mod, O[j]);
   R = R.Mul(Q).isz(4 << level);</pre>
                                                                     if (d[x][s] > -eps) return true;
                                                                     int r = -1;
for (int i = 0; i < m; ++i) {</pre>
  }
  return R.isz(L);
                                                                      if (d[i][s] < eps) continue;</pre>
 static int LinearRecursion(const V &a, const V &c,
                                                                      if (r == -1 ||
    int64_t n) { // a_n = \sum c_j a_(n-j)}
                                                                       d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  const int k = (int)a.size();
                                                                     if (r == -1) return false;
  assert((int)c.size() == k + 1);
  Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
                                                                    pivot(r, s);
  fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
  C[k] = 1
  while (n) {
                                                                  VD solve(const VVD &a, const VD &b, const VD &c) {
                                                                   m = b.size(), n = c.size();
   if (n % 2) W = W.Mul(M).DivMod(C).second;
   n /= 2, M = M.Mul(M).DivMod(C).second;
                                                                    d = VVD(m + 2, VD(n + 2))
                                                                    for (int i = 0; i < m; ++i)</pre>
                                                                    for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
  int ret = 0;
  fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
                                                                    p.resize(m), q.resize(n + 1);
                                                                    for (int i = 0; i < m; ++i) 
p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
  return ret:
                                                                    for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
                                                                    q[n] = -1, d[m + 1][n] = 1;
#undef fi
using Poly_t = Poly<998244353, 3, 1 << 20>;
                                                                    int r = 0;
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
                                                                    for (int i = 1; i < m; ++i)</pre>
                                                                     if (d[i][n + 1] < d[r][n + 1]) r = i;
5.20 Quadratic residue
                                                                    if (d[r][n + 1] < -eps) {</pre>
struct S {
                                                                     pivot(r, n);
 int MOD, w;
                                                                     if (!phase(1) || d[m + 1][n + 1] < -eps)
 int64_t x, y;
                                                                      return VD(n, -inf);
                                                                     for (int i = 0; i < m; ++i) if (p[i] == -1) {
 S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
                                                                      int s = min_element(d[i].begin(), d[i].end() - 1)
  : MOD(m), w(w_{-}), x(x_{-}), y(y_{-}) {}
 S operator*(const S &rhs) const {
                                                                           - d[i].begin();
  int w_{-} = w;
                                                                      pivot(i, s);
  if (w<sub>_</sub> == -1) w<sub>_</sub> = rhs.w;
  assert(w_! = -1 \text{ and } w_ == rhs.w);
  return { MOD, w_,
                                                                    if (!phase(0)) return VD(n, inf);
   (x * rhs.x + y * rhs.y % MOD * w) % MOD,
(x * rhs.y + y * rhs.x) % MOD };
                                                                    VD x(n);
                                                                    for (int i = 0; i < m; ++i)</pre>
                                                                     if (p[i] < n) x[p[i]] = d[i][n + 1];
                                                                    return x;
}:
                                                                  }}
int get_root(int n, int P) {
  if (P == 2 or n == 0) return n;
                                                                  5.22 Simplex Construction
  if (qpow(n, (P - 1) / 2, P) != 1) return -1;
  auto check = [&](int x) {
                                                                  Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that for all 1~\leq~j~\leq~m,
                                                                  \sum_{1 \le i \le n} A_{ji} x_i \le b_j and x_i \ge 0 for all 1 \le i \le n.
    return qpow(x, (P - 1) / 2, P); };
  if (check(n) == P-1) return -1
  int64_t a; int w; mt19937 rnd(7122);
                                                                     1. In case of minimization, let c_i' = -c_i
  do { a = rnd() % P;
w = ((a * a - n) % P + P) % P;
                                                                     2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
  } while (check(w) != P - 1);
                                                                     3. \sum_{1 < i < n} A_{ji} x_i = b_j
  return qpow(S(P, w, a, 1), (P + 1) / 2).x;
                                                                           • \sum_{1 \le i \le n} A_{ji} x_i \le b_j
5.21 Simplex
                                                                           • \sum_{1 \le i \le n} A_{ji} x_i \ge b_j
namespace simplex {
```

4. If x_i has no lower bound, replace x_i with $x_i - x_i'$

Geometry

6.1 Basic Geometry

```
#define IM imag
#define RE real
using lld = int64_t;
using llf = long double;
using PT = std::complex<lld>;
using PTF = std::complex<llf>;
auto toPTF(PT p) { return PTF{RE(p), IM(p)}; }
int sgn(11d x) { return (x > 0) - (x < 0); }</pre>
11d dot(PT a, PT b) { return RE(conj(a) * b); }
11d cross(PT a, PT b) { return IM(conj(a) * b); }
int ori(PT a, PT b, PT c) {
return sgn(cross(b - a, c - a));
bool operator<(const PT &a, const PT &b) {
return RE(a) != RE(b) ? RE(a) < RE(b) : IM(a) < IM(b);
int quad(PT p) {
 return (IM(p) == 0) // use sgn for PTF
  ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
int argCmp(PT a, PT b) {
 // -1 / 0 / 1 <-> < / == / > (atan2)
 int qa = quad(a), qb = quad(b);
 if (qa != qb) return sgn(qa - qb);
 return sgn(cross(b, a));
template <typename V> llf area(const V & pt) {
11d ret = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
 ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 return ret / 2.0;
PT rot90(PT p) { return PT{-IM(p), RE(p)}; }
PTF project(PTF p, PTF q) { // p onto q
return dot(p, q) * q / dot(q, q);
11f FMOD(11f x) {
if (x < -PI) x += PI * 2;
 if (x > PI) x -= PI * 2;
 return x;
```

6.2 Segment & Line Intersection

```
struct Segment { // closed segment
PT st, dir; // represent st + t*dir for 0 <= t <= 1
Segment(PT s, PT e) : st(s), dir(e - s) {}
static bool valid(1ld p, 1ld q) {
 // is there t s.t. 0 <= t <= 1 && qt == p ?
 if (q < 0) q = -q, p = -p;
 return 0 <= p && p <= q;
vector<PT> ends() const { return { st, st + dir }; }
template <typename T> bool isInter(T A, PT P) {
if (A.dir == PT(0)) return P == A.st; // BE CAREFUL
return cross(P - A.st, A.dir) == 0 &&
 T::valid(dot(P - A.st, A.dir), norm(A.dir));
template <typename U, typename V>
bool isInter(U A, V B) {
if (cross(A.dir, B.dir) == 0) { // BE CAREFUL
 bool res = false
 for (PT P: A.ends()) res |= isInter(B, P);
 for (PT P: B.ends()) res |= isInter(A, P);
 return res;
PT D = B.st - A.st;
11d C = cross(A.dir, B.dir);
return U::valid(cross(D, B.dir), C) &&
 V::valid(cross(D, A.dir), C);
struct Line {
PT st, ed, dir;
Line (PT s, PT e)
  : st(s), ed(e), dir(e - s) {}
PTF intersect(const Line &A, const Line &B) {
11f t = cross(B.st - A.st, B.dir) /
llf(cross(A.dir, B.dir));
```

```
18
 return toPTF(A.st) + PTF(t) * toPTF(A.dir);
6.3 2D Convex Hull
void make_hull(vector<pll> &dots) { // n=1 => ans = {}
 sort(dots.begin(), dots.end());
 vector<pll> ans(1, dots[0]);
 for (int ct = 0; ct < 2; ++ct, reverse(ALL(dots)))</pre>
  for (int i = 1, t = SZ(ans); i < SZ(dots); i++) {</pre>
   while (SZ(ans) > t && ori(
     ans[SZ(ans) - 2], ans.back(), dots[i]) <= 0)
    ans.pop_back();
   ans.pb(dots[i]);
 ans.pop_back(), ans.swap(dots);
6.4 3D Convex Hull
// return the faces with pt indexes
int flag[MXN][MXN];
struct Point{
 ld x,y,z;
 Point operator * (const ld &b) const {
 return (Point) {x*b,y*b,z*b}; }
Point operator * (const Point &b) const {
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
Point ver(Point a, Point b, Point c) {
  return (b - a) * (c - a);}
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0;
 REP(i,n) REP(j,n) flag[i][j] = 0;
 vector<Face> now;
 now.emplace_back(0,1,2);
 now.emplace_back(2,1,0)
 for (int i=3; i<n; i++){
  ftop++; vector<Face> next;
  REP(j, SZ(now)) {
  Face& f=now[j]; int ff = 0;
   ld d=(pt[i]-pt[f.a]).dot(
     ver(pt[f.a], pt[f.b], pt[f.c]));
   if (d <= 0) next.push_back(f);</pre>
   if (d > 0) ff=ftop;
   else if (d < 0) ff=-ftop</pre>
   flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
  REP(j, SZ(now)) {
  Face& f=now[j];
   if (flag[f.a][f.b] > 0 &&
     flag[f.a][f.b] != flag[f.b][f.a])
    next.emplace_back(f.a,f.b,i);
   if (flag[f.b][f.c] > 0 &&
     flag[f.b][f.c] != flag[f.c][f.b])
    next.emplace_back(f.b,f.c,i);
   if (flag[f.c][f.a] > 0 &&
  flag[f.c][f.a] != flag[f.a][f.c])
    next.emplace_back(f.c,f.a,i);
  now=next;
 return now;
6.5 2D Farthest Pair
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++)</pre>
 while(abs(cross(stk[i+1]-stk[i],
   stk[(pos+1)%n]-stk[i])) >
   abs(cross(stk[i+1]-stk[i],
   stk[pos]-stk[i]))) pos = (pos+1)%n;
 ans = max({ans, dis(stk[i], stk[pos]),
    dis(stk[i+1], stk[pos])});
```

6.6 kD Closest Pair (3D ver.)

shuffle(v.begin(), v.end(), mt19937()); unordered_map<lld, unordered_map<lld,</pre>

11f solve(vector<P> v) {

while (pt.size() > 0 \

&& ori(R.st, R.ed, pt.front()) < 0) \

pt.pop_front(), que.pop_front();

```
unordered_map<lld, int>>> m;
                                                                 POP(lines[i], lines[i]);
                                                                 pt.push_back(intersect(que.back(), lines[i]));
llf d = dis(v[0], v[1]);
auto Idx = [\&d] (11f x) -> 11d {
                                                                 que.push_back(lines[i]);
  return round(x * 2 / d) + 0.1;
auto rebuild_m = [&m, &v, &Idx](int k) {
                                                               POP(que.front(), que.back())
  m.clear();
                                                               if (que.size() <= 1 ||</pre>
                                                                 argCmp(que.front().dir, que.back().dir) == 0)
  for (int i = 0; i < k; ++i)
  m[Idx(v[i].x)][Idx(v[i].y)]
                                                                 return 0:
    [Idx(v[i].z)] = i;
                                                               pt.push_back(intersect(que.front(), que.back()));
 }; rebuild_m(2);
                                                               return area(pt);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
 const 11d kx = Idx(v[i].x), ky = Idx(v[i].y),
                                                             6.9 Minkowski Sum
     kz = Idx(v[i].z); bool found = false;
  for (int dx = -2; dx <= 2; ++dx) {
                                                            vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
   const 11d nx = dx + kx;
                                                              hull(A), hull(B);
                                                              vector<pl1> C(1, A[0] + B[0]), s1, s2;
   if (m.find(nx) == m.end()) continue;
   auto& mm = m[nx];
                                                              for(int i = 0; i < SZ(A); ++i)
s1.pb(A[(i + 1) % SZ(A)] - A[i]);</pre>
   for (int dy = -2; dy <= 2; ++dy) {
    const 11d ny = dy + ky;
                                                              for(int i = 0; i < SZ(B); i++)
    if (mm.find(ny) == mm.end()) continue;
                                                               s2.pb(B[(i + 1) % SZ(B)] - B[i]);
    auto& mmm = mm[ny];
                                                              for(int p1 = 0, p2 = 0; p1 < SZ(A) || p2 < SZ(B);)</pre>
    for (int dz = -2; dz <= 2; ++dz) {
                                                               if (p2 >= SZ(B)
     const 11d nz = dz + kz;
                                                                 || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
     if (mmm.find(nz) == mmm.end()) continue;
                                                                C.pb(C.back() + s1[p1++]);
     const int p = mmm[nz];
     if (dis(v[p], v[i]) < d) {
  d = dis(v[p], v[i]);</pre>
                                                                C.pb(C.back() + s2[p2++]);
                                                              return hull(C), C;
      found = true;
     }
                                                             6.10 Circle Class
                                                            struct Circle { PTF o; llf r; };
  if (found) rebuild_m(i + 1);
                                                             vector<llf> intersectAngle(Circle A, Circle B) {
 else m[kx][ky][kz] = i;
                                                              PTF dir = B.o - A.o; llf d2 = norm(dir);
                                                              if (norm(A.r - B.r) >= d2) // norm(x) := |x|^2
return d;
                                                               if (A.r < B.r) return {-PI, PI}; // A in B</pre>
                                                               else return {}; // B in A
                                                              if (norm(A.r + B.r) <= d2) return {};</pre>
6.7 Simulated Annealing
                                                              11f dis = abs(dir), theta = arg(dir);
11f anneal() {
                                                              11f phi = acos((A.r * A.r + d2 - B.r * B.r) /
mt19937 rnd_engine( seed );
                                                                (2 * A.r * dis));
uniform_real_distribution< llf > rnd( 0, 1 );
                                                              11f L = FMOD(theta - phi), R = FMOD(theta + phi);
const llf dT = 0.001;
                                                              return { L, R };
 // Argument p
11f S_cur = calc( p ), S_best = S_cur;
for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
                                                             vector<PTF> intersectPoint(Circle a, Circle b) {
 // Modify p to p_prime
                                                              llf d = abs(a.o - b.o);
 const llf S_prime = calc( p_prime );
                                                              if (d >= b.r+a.r || d <= abs(b.r-a.r)) return {};</pre>
  const 11f delta_c = S_prime - S_cur
                                                              11f dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
 llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
                                                              PTF dir = (a.o - b.o) / d;
 if ( rnd( rnd_engine ) <= prob )</pre>
                                                              PTF u = dir*d1 + b.o;
   S_cur = S_prime, p = p_prime;
                                                              PTF v = rot90(dir) * sqrt(max<llf>(0, b.r*b.r-d1*d1));
 if ( S_prime < S_best ) // find min</pre>
                                                              return {u + v, u - v};
   S_best = S_prime, p_best = p_prime;
return S_best;
                                                             6.11 Intersection of line and Circle
                                                            vector<PTF> line_interCircle(const PTF &p1,
6.8 Half Plane Intersection
                                                               const PTF &p2, const PTF &c, const double r)
                                                              PTF ft = p1 + project(c-p1, p2-p1), vec = p2-p1;
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
                                                              llf dis = abs(c - ft)
bool operator<(const Line &lhs, const Line &rhs) {</pre>
                                                              if (abs(dis - r) < eps) return {ft};</pre>
  if (int cmp = argCmp(lhs.dir, rhs.dir))
                                                              if (dis > r) return {};
    return cmp == -1;
                                                              vec = vec * sqrt(r * r - dis * dis) / abs(vec);
  return ori(lhs.st, lhs.ed, rhs.st) < 0;</pre>
                                                              return {ft + vec, ft - vec};
// intersect function is in "Segment Intersect"
                                                             6.12 Intersection of Polygon and Circle
11f HPI(vector<Line> &lines) {
                                                            // Divides into multiple triangle, and sum up
  sort(lines.begin(), lines.end());
                                                             // test by HDU2892
  deque<Line> que;
 deque<PTF> pt;
                                                            11f _area(PTF pa, PTF pb, llf r)
                                                              if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
  que.push_back(lines[0]);
  for (int i = 1; i < (int)lines.size(); i++) {</pre>
                                                              if (abs(pb) < eps) return 0;</pre>
                                                              11f S, h, theta;
    if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
                                                              llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
#define POP(L, R) \
                                                              11f cosB = dot(pb, pb - pa) / a / c, B = acos(cosB);
                                                              llf cosC = dot(pa, pb) / a / b, C = acos(cosC);
    while (pt.size() > 0 \
      && ori(L.st, L.ed, pt.back()) < 0) \
                                                              if (a > r) ·
      pt.pop_back(), que.pop_back(); \
                                                              S = (C / 2) * r * r;
```

h = a * b * sin(C) / c

if (h < r && B < PI / 2)

S = (acos(h / r) * r * r - h * sqrt(r*r - h*h));

} else if (b > r) {

```
theta = PI - B - asin(sin(B) / r * a);
 S = 0.5 * a*r * sin(theta) + (C - theta) / 2 * r*r;
 S = 0.5 * sin(C) * a * b;
return S;
11f area_poly_circle(const vector<PTF> &poly,
 const PTF &0, const llf r) {
 11f S = 0;
for (int i = 0, N = poly.size(); i < N; ++i)</pre>
 S += _area(poly[i] - 0, poly[(i + 1) % N] - 0, r) *
    ori(0, poly[i], poly[(i + 1) % N]);
 return fabs(S);
6.13 Point & Hulls Tangent
#define above(P, Vi, Vj) (ori(P, Vi, Vj) > 0) // true
    if Vi is above Vj
#define below(P, Vi, Vj) (ori(P, Vi, Vj) < 0) // true</pre>
    if Vi is below Vj
// Rtangent_PointPolyC(): binary search for convex
    polygon right tangent
   Input: P = a \ 2D \ point \ (exterior \ to \ the \ polygon)
        n = number of polygon vertices
        V = array of vertices for a 2D convex polygon
11
    with V[n] = V[0]
// Return: index "i" of rightmost tangent point V[i]
int Rtangent_PointPolyC(PT P, int n, PT *V) {
int a, b, c;
int upA, dnC;
if (below(P, V[1], V[0]) && !above(P, V[n - 1], V[0]))
 return 0;
for (a = 0, b = n;;) {
 c = (a + b) / 2
  dnC = below(P, V[c + 1], V[c]);
  if (dnC && !above(P, V[c - 1], V[c]))
   return c;
 upA = above(P, V[a + 1], V[a]);
  if (upA) {
  if (dnC) {
   b = c;
   } else {
    if (above(P, V[a], V[c]))
     b = c;
    else
     a = c:
   }
  } else {
   if (!dnC) {
   a = c;
   } else {
    if (below(P, V[a], V[c]))
    b = c;
   else
     a = c;
  }
}
// Ltangent_PointPolyC(): binary search for convex
    polygon left tangent
    Input: P = a \ 2D \ point \ (exterior \ to \ the \ polygon)
        n = number of polygon vertices
        V = array of vertices for a 2D convex polygon
//
    with V[n]=V[0]
   Return: index "i" of leftmost tangent point V[i]
int Ltangent_PointPolyC(PT P, int n, PT *V) {
int a, b, c;
int dnA, dnC;
if (above(P, V[n - 1], V[0]) && !below(P, V[1], V[0]))
 return 0;
for (a = 0, b = n;;) {
 c = (a + b) / 2;
dnC = below(P, V[c + 1], V[c]);
 if (above(P, V[c - 1], V[c]) && !dnC)
```

```
dnA = below(P, V[a + 1], V[a]);
 if (dnA)
  if (!dnC) {
  b = c;
  } else {
   if (below(P, V[a], V[c]))
   b = c;
   else
    a = c:
 } else {
  if (dnC) {
   a = c;
  } else {
   if (above(P, V[a], V[c]))
   b = c;
   else
    a = c:
}
```

6.14 Convex Hulls Tangent

```
// RLtangent_PolyPolyC(): get the RL tangent between
    two convex polygons
   Input: m = number of vertices in polygon 1
       V = array of vertices for convex polygon 1 with
//
     V[m]=V[0]
//
       n = number of vertices in polygon 2
       W = array of vertices for convex polygon 2 with
//
     W[n]=W[0]
   Output: *t1 = index of tangent point V[t1] for
//
    polygon 1
        *t2 = index of tangent point W[t2] for polygon
void RLtangent_PolyPolyC(int m, PT *V, int n, PT *W,
    int *t1, int *t2) {
 int ix1, ix2; // search indices for polygons 1 and 2
 // first get the initial vertex on each polygon
 ix1 = Rtangent_PointPolyC(W[0], m, V); // right
    tangent from W[0] to V
 ix2 = Ltangent_PointPolyC(V[ix1], n, W); // left
    tangent from V[ix1] to W
 // ping-pong linear search until it stabilizes
 int done = false; // flag when done
 while (done == false) {
 done = true; // assume done until..
  while (ori(W[ix2], V[ix1], V[ix1 + 1]) <= 0) {</pre>
   ++ix1; // get Rtangent from W[ix2] to V
 done = false; // not done if had to adjust this
 }
 *t1 = ix1;
 *t2 = ix2:
 return;
```

6.15 Tangent line of Two Circle

```
vector<Line>
tanline(const Circle &c1, const Circle &c2, int sign1){
 // sign1 = 1 for outer tang, -1 for inter tang
 vector<Line> ret:
 if (norm(c1.o - c2.o) < eps) return ret;</pre>
 11f d = abs(c1.o - c2.o);
 PTF v = (c2.o - c1.o) / d;
 11f c = (c1.r - sign1 * c2.r) / d;
 if (c * c > 1) return ret;
 llf h = sqrt(max<llf>(0, 1 - c * c));
 for (int sign2 : {1, -1}) {
 PTF n = c * v + sign2 * h * rot90(v);
 PTF p1 = c1.o + n * c1.r;
  PTF p2 = c2.0 + n * (c2.r * sign1);
  if (norm(p2 - p1) < eps)
  p2 = p1 + rot90(c2.o - c1.o);
  ret.push_back({p1, p2});
```

```
return tree+M;
return ret;
}
                                                                  int touch(Node* r, int x, int y, LL d2){
                                                                   LL dis = sqrt(d2)+1;
6.16 Minimum Covering Circle
                                                                   if (x<r->x1-dis || x>r->x2+dis ||
template<typename P>
                                                                     y<r->y1-dis || y>r->y2+dis)
Circle getCircum(const P &a, const P &b, const P &c){
                                                                    return 0;
                                                                   return 1;
 Real a1 = a.x-b.x, b1 = a.y-b.y;
 Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
 Real a2 = a.x-c.x, b2 = a.y-c.y;
                                                                  void nearest(Node* r,int x,int y,int &mID,LL &md2) {
 Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
                                                                   if (!r || !touch(r, x, y, md2)) return;
                                                                   LL d2 = dis2(r->x, r->y, x, y);
 cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
                                                                   if (d2 < md2 | | (d2 == md2 && mID < r->id)) {
 cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2)
                                                                    mID = r->id;
                                                                    md2 = d2;
 cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
 return cc;
                                                                   // search order depends on split dim
                                                                   if ((r->f == 0 && x < r->x) ||
                                                                     (r->f == 1 \&\& y < r->y)) {
template<typename P>
                                                                    nearest(r->L, x, y, mID, md2);
nearest(r->R, x, y, mID, md2);
Circle MinCircleCover(const vector<P>& pts){
 random_shuffle(pts.begin(), pts.end());
 Circle c = { pts[0], 0 };
                                                                   } else {
 for(int i=0;i<(int)pts.size();i++){</pre>
                                                                    nearest(r->R, x, y, mID, md2);
 if (dist(pts[i], c.o) <= c.r) continue;</pre>
                                                                    nearest(r->L, x, y, mID, md2);
  c = { pts[i], 0 };
  for (int j = 0; j < i; j++) {
   if(dist(pts[j], c.o) <= c.r) continue;</pre>
                                                                  int query(int x, int y) {
   c.o = (pts[i] + pts[j]) / 2;
                                                                   int id = 1029384756;
   c.r = dist(pts[i], c.o);
                                                                   LL d2 = 102938475612345678LL;
   for (int k = 0; k < j; k++) {
                                                                   nearest(root, x, y, id, d2);
    if (dist(pts[k], c.o) <= c.r) continue;</pre>
                                                                   return id:
    c = getCircum(pts[i], pts[j], pts[k]);
                                                                 } tree;
  }
                                                                 6.18 Rotating Sweep Line
                                                                 void rotatingSweepLine(pair<int, int> a[], int n) {
 return c;
                                                                  vector<pair<int, int>> 1;
                                                                  1.reserve(n * (n - 1) / 2)
        KDTree (Nearest Point)
6.17
                                                                  for (int i = 0; i < n; ++i)
for (int j = i + 1; j < n; ++j)
const int MXN = 100005;
struct KDTree {
                                                                    1.emplace_back(i, j);
                                                                  sort(1.begin(), 1.end(), [&a](auto &u, auto &v){
 struct Node {
                                                                   1ld udx = a[u.first].first - a[u.second].first;
  int x,y,x1,y1,x2,y2;
 int id,f;
Node *L, *R;
                                                                   11d udy = a[u.first].second - a[u.second].second;
                                                                   1ld vdx = a[v.first].first - a[v.second].first;
1ld vdy = a[v.first].second - a[v.second].second;
 } tree[MXN], *root;
                                                                   if (udx == 0 or vdx == 0) return not udx == 0;
                                                                   int s = sgn(udx * vdx);
LL dis2(int x1, int y1, int x2, int y2) {
LL dx = x1-x2, dy = y1-y2;
                                                                   return udy * vdx * s < vdy * udx * s;
 return dx*dx+dy*dy;
                                                                  }):
                                                                  vector<int> idx(n), p(n);
                                                                  iota(idx.begin(), idx.end(), 0);
sort(idx.begin(), idx.end(), [&a](int i, int j){
 static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
 static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 void init(vector<pair<int,int>> ip) {
                                                                   return a[i] < a[j]; });</pre>
                                                                  for (int i = 0; i < n; ++i) p[idx[i]] = i;
for (auto [i, j]: 1) {</pre>
 n = ip.size();
  for (int i=0; i<n; i++) {</pre>
   tree[i].id = i;
                                                                   // do here
   tree[i].x = ip[i].first;
                                                                   swap(p[i], p[j]);
   tree[i].y = ip[i].second;
                                                                   idx[p[i]] = i, idx[p[j]] = j;
  root = build_tree(0, n-1, 0);
                                                                 6.19
                                                                        Circle Cover
 Node* build_tree(int L, int R, int d) {
                                                                 const int N = 1021;
  if (L>R) return nullptr
  int M = (L+R)/2; tree[M].f = d%2;
                                                                 struct CircleCover {
  nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
                                                                  int C;
  tree[M].x1 = tree[M].x2 = tree[M].x;
                                                                  Cir c[N]
                                                                  bool g[N][N], overlap[N][N];
// Area[i] : area covered by at least i circles
  tree[M].y1 = tree[M].y2 = tree[M].y;
  tree[M].L = build_tree(L, M-1, d+1);
                                                                  double Area[ N ];
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
                                                                  void init(int _C){ C = _C;}
                                                                  struct Teve {
   tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
                                                                   PTF p; double ang; int add;
                                                                   Teve() {}
   tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
                                                                   Teve(PTF _a, double _b, int _c):p(_a), ang(_b), add(
  tree[M].R = build_tree(M+1, R, d+1);
  if (tree[M].R) {
                                                                   bool operator<(const Teve &a)const
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
                                                                   {return ang < a.ang;}
   tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
                                                                  }eve[N * 2];
   tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
                                                                  // strict: x = 0, otherwise x = -1
                                                                  bool disjuct(Cir &a, Cir &b, int x)
```

{return sign(abs(a.0 - b.0) - a.R - b.R) > x;}

```
bool contain(Cir &a, Cir &b, int x)
{return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
bool contain(int i, int j) {
 /* c[j] is non-strictly in c[i]. */
return (sign(c[i].R - c[j].R) > 0 || (sign(c[i].R - c
    [j].R) == 0 && i < j)) && contain(c[i], c[j], -1);
void solve(){
 fill_n(Area, C + 2, 0);
 for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)
   overlap[i][j] = contain(i, j);
 for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j</pre>
   g[i][j] = !(overlap[i][j] || overlap[j][i] ||
      disjuct(c[i], c[j], -1));
 for(int i = 0; i < C; ++i){</pre>
  int E = 0, cnt = 1;
  for(int j = 0; j < C; ++j)
   if(j != i && overlap[j][i])
    ++cnt:
  for(int j = 0; j < C; ++j)
   if(i != j && g[i][j]) {
    auto IP = intersectPoint(c[i], c[j]);
    PTF aa = IP[0], bb = IP[1];
    llf A = arg(aa-c[i].0), B = arg(bb-c[i].0);
eve[E++] = Teve(bb,B,1), eve[E++]=Teve(aa,A,-1);
    if(B > A) ++cnt;
  if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
   sort(eve, eve + E);
   eve[E] = eve[0];
   for(int j = 0; j < E; ++j){
    cnt += eve[j].add;
    Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
    double theta = eve[j + 1].ang - eve[j].ang;
    if (theta < 0) theta += 2. * pi;</pre>
    Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
 }
}
```

7 Stringology

7.1 Hash

```
class Hash {
 private:
  static constexpr int P = 127, Q = 1051762951;
  vector<int> h, p;
  void init(const string &s){
   h.assign(s.size()+1, 0); p.resize(s.size()+1);
for (size_t i = 0; i < s.size(); ++i)
     h[i + 1] = add(mul(h[i], P), s[i]);
   generate(p.begin(), p.end(), [x=1, y=1, this]()
mutable{y=x;x=mul(x,P);return y;});
  int query(int 1, int r){ // 1-base (1, r]
  return sub(h[r], mul(h[1], p[r-1]));}
```

7.2 Suffix Array

```
namespace sfx {
bool _t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], _c[maxn * 2];
int x[maxn], _p[maxn], _q[maxn * 2];
// sa[i]: sa[i]-th suffix is the
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *a, int *c, int n, int z) {
 memset(a, 0, sizeof(int) * n);
 memcpy(x, c, sizeof(int) * z);
void induce(int *a,int *c,int *s,bool *t,int n,int z){
memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
```

```
if (a[i] && !t[a[i] - 1])
   a[x[s[a[i] - 1]]++] = a[i] - 1;
 memcpy(x, c, sizeof(int) * z);
 for (int i = n - 1; i >= 0; --i)
if (a[i] && t[a[i] - 1])
   a[--x[s[a[i] - 1]]] = a[i] - 1;
void sais(int *s, int *a, int *p, int *q,
 bool *t, int *c, int n, int z) {
bool uniq = t[n - 1] = true;
 int nn=0, nmxz=-1, *nsa = a+n, *ns=s+n, last=-1;
 memset(c, 0, sizeof(int) * z);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
 if (uniq) {
  for (int i = 0; i < n; ++i) a[--c[s[i]]] = i;
  return:
 for (int i = n - 2; i >= 0; --i)
  t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 pre(a, c, n, z);
 for (int i = 1; i <= n - 1; ++i)</pre>
  if (t[i] && !t[i - 1])
   a[--x[s[i]]] = p[q[i] = nn++] = i;
 induce(a, c, s, t, n, z);
for (int i = 0; i < n; ++i)
  if (a[i] && t[a[i]] && !t[a[i] - 1]) {
  bool neq = last < 0 || \</pre>
   memcmp(s + a[i], s + last,
(p[q[a[i]] + 1] - a[i]) * sizeof(int));
  ns[q[last = a[i]]] = nmxz += neq;
 sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
 pre(a, c, n, z);
 for (int i = nn - 1; i >= 0; --i)
  a[--x[s[p[nsa[i]]]]] = p[nsa[i]];
 induce(a, c, s, t, n, z);
void build(const string &s) {
 const int n = int(s.size());
 for (int i = 0; i < n; ++i) _s[i] = s[i];
 _s[n] = 0; // s shouldn't contain 0
 sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
 for(int i = 0; i < n; ++i) rev[sa[i] = sa[i+1]] = i;</pre>
 int ind = hi[0] = 0;
 for (int i = 0; i < n; ++i) {
  if (!rev[i]) {
   ind = 0:
   continue;
  while (i + ind < n && \</pre>
   s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  hi[rev[i]] = ind ? ind-- : 0;
}}
```

7.3 Suffix Automaton

```
struct SuffixAutomaton {
 struct node {
  int ch[K], len, fail, cnt, indeg;
  node(int L = 0) : ch{}, len(L), fail(0), cnt(0),
    indeg(0) {}
 } st[N];
 int root, last, tot;
 void extend(int c) {
  int cur = ++tot;
  st[cur] = node(st[last].len + 1);
  while (last && !st[last].ch[c]) {
    st[last].ch[c] = cur;
    last = st[last].fail;
  if (!last) {
    st[cur].fail = root;
  } else {
    int q = st[last].ch[c];
    if (st[q].len == st[last].len + 1) {
      st[cur].fail = q;
    } else {
      int clone = ++tot;
      st[clone] = st[q];
      st[clone].len = st[last].len + 1;
      st[st[cur].fail = st[q].fail = clone].cnt = 0;
```

```
while (last && st[last].ch[c] == q) {
                                                                 for(char c: s) t += c, t += '.';
                                                                int 1 = 0, r = 0, ans = 0;
        st[last].ch[c] = clone;
                                                                for (int i = 1; i < t.length(); ++i) {
z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
while (i - z[i] >= 0 && i + z[i] < t.length()) {</pre>
        last = st[last].fail;
    }
  }
                                                                   if(t[i - z[i]] == t[i + z[i]]) ++z[i];
  st[last = cur].cnt += 1;
                                                                   else break:
 void init(const char* s) {
                                                                 if (i + z[i] > r) r = i + z[i], l = i;
 root = last = tot = 1;
  st[root] = node(0);
                                                                for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
  for (char c; c = *s; ++s) extend(c - 'a');
 int q[N];
                                                               7.7
                                                                      Lexico Smallest Rotation
 void dp() {
 for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
                                                               string mcp(string s) {
                                                                int n = s.length();
  int head = 0, tail = 0;
                                                                s += s; int i = 0, j = 1;
  for (int i = 1; i <= tot; i++)</pre>
                                                                while (i < n && j < n) {</pre>
    if (st[i].indeg == 0) q[tail++] = i;
                                                                 int k = 0;
  while (head != tail) {
                                                                 while (k < n \&\& s[i + k] == s[j + k]) k++;
    int now = q[head++];
if (int f = st[now].fail) {
                                                                 ((s[i+k] \le s[j+k]) ? j : i) += k + 1;
                                                                 j += (i == j);
      st[f].cnt += st[now].cnt;
      if (--st[f].indeg == 0) q[tail++] = f;
                                                                return s.substr(i < n ? i : j, n);</pre>
  }
                                                               7.8 Main Lorentz
 int run(const char* s) {
                                                               vector<tuple<tuple<size_t, size_t, int, int>>> reps;
  int now = root;
                                                               void find_repetitions(const string &s, int shift = 0) {
  for (char c; c = *s; ++s) {
                                                                if (s.size() <= 1)
    if (!st[now].ch[c -= 'a']) return 0;
    now = st[now].ch[c];
                                                                const size_t nu = s.size() / 2, nv = s.size() - nu;
                                                                string u = s.substr(0, nu), v = s.substr(nu);
  return st[now].cnt;
                                                                string ru(u.rbegin(), u.rend());
 }
                                                                string rv(v.rbegin(), v.rend());
} SAM;
                                                                find_repetitions(u, shift);
                                                                find_repetitions(v, shift + nu);
7.4 KMP
                                                                auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u),
z3 = Zalgo(ru + '#' + rv), z4 = Zalgo(v);
vector<int> kmp(const string &s) {
 vector<int> f(s.size(), θ);
                                                                for (size_t cntr = 0; cntr < s.size(); cntr++) {</pre>
 /* f[i] = length of the longest prefix
                                                                 size_t 1; int k1, k2;
   (excluding s[0:i]) such that it coincides
                                                                 if (cntr < nu) {</pre>
   with the suffix of s[0:i] of the same length */
                                                                   1 = nu - cntr
 /* i + 1 - f[i] is the length of the
                                                                   k1 = 1 < z1.size() ? z1[1] : 0;
   smallest recurring period of s[0:i] */
                                                                   k2 = n + 1 - 1 < z2.size() ? z2[n + 1 - 1] : 0;
 int k = 0;
                                                                 } else {
 for (int i = 1; i < (int)s.size(); ++i) {</pre>
                                                                   1 = cntr - nu + 1;
  while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
                                                                   k1 = n + 1 - 1 < z3.size() ? z3[n + 1 - 1] : 0;
  if (s[i] == s[k]) ++k;
                                                                  k2 = 1 < z4.size() ? z4[1] : 0;
 f[i] = k;
                                                                 if (k1 + k2 >= 1)
return f;
                                                                   reps.emplace_back(cntr, 1, k1, k2);
vector<int> search(const string &s, const string &t) {
 // return 0-indexed occurrence of t in s
 vector<int> f = kmp(t), r;
for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
                                                               7.9 BWT
                                                               struct BurrowsWheeler{
  while(k > 0 && (k==(int)t.size() || s[i]!=t[k]))
                                                               #define SIGMA 26
   k = f[k - 1]
                                                               #define BASE 'a
  if (s[i] == t[k]) ++k;
                                                                vector<int> v[ SIGMA ];
  if (k == (int)t.size()) r.push_back(i-t.size()+1);
                                                                void BWT(char* ori, char* res){
                                                                 // make ori -> ori + ori
 return res:
                                                                 // then build suffix array
7.5 Z value
                                                                void iBWT(char* ori, char* res){
                                                                 for( int i = 0 ; i < SIGMA ; i ++ )</pre>
vector<int> Zalgo(const string &s) {
                                                                  v[ i ].clear()
 vector<int> z(s.size(), s.size());
                                                                 int len = strlen( ori );
 for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
                                                                 for( int i = 0 ; i < len ; i ++ )</pre>
  int j = clamp(r - i, 0, z[i - 1]);
                                                                  v[ ori[i] - BASE ].push_back( i );
  for (; i + j < z[0] \text{ and } s[i + j] == s[j]; ++j);
                                                                 vector<int> a;
  if (i + (z[i] = j) > r) r = i + z[1 = i];
                                                                 for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
 }
                                                                   for( auto j : v[ i ] ){
 return z;
                                                                    a.push_back( j );
                                                                    ori[ ptr ++ ] = BASE + i;
7.6 Manacher
                                                                 for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
int z[maxn];
int manacher(const string& s) {
  string t = ".";
                                                                  ptr = a[ ptr ];
```

```
res[ len ] = 0;
}
} bwt;
```

7.10 Palindromic Tree

```
struct palindromic_tree{
struct node{
  int next[26],f,len;
  int cnt,num,st,ed; // num = depth of fail link
  node(\textbf{int }1\text{=}0)\text{:}f(0)\text{,len}(1)\text{,cnt}(0)\text{,num}(0)\text{ }\{
   memset(next, 0, sizeof(next)); }
}:
vector<node> st;
vector<char> s;
int last.n;
void init(){
 st.clear();s.clear();last=1; n=0;
 st.push_back(0);st.push_back(-1);
  st[0].f=1;s.push_back(-1); }
 int getFail(int x){
 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
  return x;}
void add(int c){
  s.push_back(c-='a'); ++n;
  int cur=getFail(last);
  if(!st[cur].next[c]){
   int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
  st[now].num=st[st[now].f].num+1;
  last=st[cur].next[c];
  ++st[last].cnt;}
 void dpcnt() { // cnt = #occurence in whole str
  for (int i=st.size()-1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
int size(){ return st.size()-2;}
} pt;
int main() {
string s; cin >> s; pt.init();
for (int i=0; i<SZ(s); i++)</pre>
 int prvsz = pt.size(); pt.add(s[i]);
 if (prvsz != pt.size()) {
   int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [1,r]: s.substr(1, r-1+1)
 }
return 0;
```

8 Misc

8.1 Theorems

8.1.1 Sherman-Morrison formula

$$(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1+v^{\mathsf{T}}A^{-1}u}$$

8.1.2 Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i), L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

8.1.3 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij}=x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij}=-d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.4 Cayley's Formula

- Given a degree sequence d_1,d_2,\ldots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.5 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \leq k \leq n$.

8.1.6 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.7 Euler's planar graph formula

```
V - E + F = C + 1, E \le 3V - 6(?)
```

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2)$, find maximum $S\in I_1\cap I_2$. For each iteration, build the directed graph and find a shortest path from s to t.

```
• s \to x: S \sqcup \{x\} \in I_1

• x \to t: S \sqcup \{x\} \in I_2

• y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1 (y is in the unique circuit of S \sqcup \{x\})

• x \to y: S \setminus \{y\} \sqcup \{x\} \in I_2 (y is in the unique circuit of S \sqcup \{x\})
```

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|.$ In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S,$ resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

8.2 Bitset LCS

```
scanf("%d%d", &n, &m), u = n / 64 + 1;
for (int i = 1, c; i <= n; i++)
  scanf("%d", &c), p[c].set(i);
for (int i = 1, c; i <= m; i++) {
  scanf("%d", &c), (g = f) |= p[c];
  f.shiftLeftByOne(), f.set(0);
  ((f = g - f) ^= g) &= g;
}
printf("%d\n", f.count());</pre>
```

8.3 Prefix Substring LCS

```
void all_lcs(string s, string t) { // 0-base
vector<int> h(SZ(t));
iota(ALL(h), 0);
for (int a = 0; a < SZ(s); ++a) {
  int v = -1;
  for (int c = 0; c < SZ(t); ++c)
   if (s[a] == t[c] || h[c] < v)
      swap(h[c], v);
  // LCS(s[0, a], t[b, c]) =
  // c - b + 1 - sum([h[i] >= b] | i <= c)
  // h[i] might become -1 !!
}
</pre>
```

8.4 Convex 1D/1D DP

```
struct segment {
 int i, 1, r
 segment() {}
 segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
 dp[0] = 0;
 deque<segment> dq; dq.push_back(segment(0, 1, n));
 for (int i = 1; i <= n; ++i) {</pre>
  dp[i] = f(dq.front().i, i);
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().1 = i + 1
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
   f(i, dq.back().1) < f(dq.back().i, dq.back().1))
    dq.pop_back();
```

```
if (dq.size()) {
 int d = 1 << 20, c = dq.back().1;</pre>
 while (d >>= 1) if (c + d <= dq.back().r)</pre>
  if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
 dq.back().r = c; seg.l = c + 1;
if (seg.1 <= n) dq.push_back(seg);</pre>
```

ConvexHull Optimization

```
struct L {
mutable int64_t a, b, p;
bool operator<(const L &r) const { return a < r.a; }</pre>
bool operator<(int64_t x) const { return p < x; }</pre>
struct DynamicHull : multiset<L, less<>> {
static const int64_t kInf = 1e18;
bool Isect(iterator x, iterator y) {
 auto Div = [](int64_t a, int64_t b) {
   return a / b - ((a ^ b) < 0 && a % b); }
 if (y == end()) { x->p = kInf; return false; }
 if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
 else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
void Insert(int64_t a, int64_t b) {
 auto z = insert({a, b, 0}), y = z++, x = y;
 while (Isect(y, z)) z = erase(z);
 if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
 while ((y = x) != begin() && (--x)->p >= y->p)
  Isect(x, erase(y));
int64_t Query(int64_t x) {
 auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
};
```

Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
int s = 0:
for (int i = 2; i <= n; i++)
 s = (s + m) \% i;
return s:
// died at kth
int kth(int n, int m, int k){
if (m == 1) return n-1;
for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
return k;
```

8.7 Tree Knapsack

```
int dp[N][K]; PII obj[N];
vector<int> G[N];
void dfs(int u, int mx){
 for(int s: G[u]) {
  if(mx < obj[s].first) continue;</pre>
  for(int i=0;i<=mx-obj[s].FF;i++)</pre>
   dp[s][i] = dp[u][i];
  dfs(s, mx - obj[s].first);
  for(int i=obj[s].FF;i<=mx;i++)</pre>
   dp[u][i] = max(dp[u][i],
    dp[s][i - obj[s].FF] + obj[s].SS);
```

8.8 N Queens Problem

```
vector< int > solve( int n ) {
// no solution when n=2, 3
vector< int > ret;
if ( n % 6 == 2 ) {
 for ( int i = 2 ; i <= n ; i += 2 )
  ret.push_back( i );
 ret.push_back( 3 ); ret.push_back( 1 );
 for ( int i = 7 ; i <= n ; i += 2 )
  ret.push_back( i );
 ret.push_back( 5 );
} else if ( n % 6 == 3 ) {
```

```
for ( int i = 4 ; i <= n ; i += 2 )
 ret.push_back( i );
 ret.push_back( 2 );
 for ( int i = 5 ; i <= n ; i += 2 )
ret.push_back( i );</pre>
 ret.push_back( 1 ); ret.push_back( 3 );
} else {
 for ( int i = 2 ; i <= n ; i += 2 )
 ret.push_back( i );
for ( int i = 1 ; i <= n ; i += 2 )
  ret.push_back( i );
return ret:
```

8.9 Stable Marriage

```
1: Initialize m \in M and w \in W to free
2: while \exists free man m who has a woman w to propose to do
         w \leftarrow \text{first woman on } m \text{'s list to whom } m \text{ has not yet proposed}
        if \exists \mbox{ some pair } (m',w) \mbox{ then }
            if w prefers m to m' then
                \dot{m'} \leftarrow \textit{free}
                 (m,w) \leftarrow \mathsf{engaged}
8:
            end if
9:
        else
10:
             (m,w) \leftarrow \mathsf{engaged}
        end if
12: end while
```

8.10 Binary Search On Fraction

```
struct Q {
 11 p, q;
 Q go(Q b, 11 d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and \theta <= p,q <= N
Q frac_bs(ll N) {
 Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
 assert(pred(hi));
 bool dir = 1, L = 1, H = 1;
 for (; L || H; dir = !dir) {
  11 len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
   if (Q mid = hi.go(lo, len + step);
     mid.p > N || mid.q > N || dir ^ pred(mid))
    t++;
   else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
 return dir ? hi : lo;
```