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	5.1 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.16 5.17 5.18 5.20 5.21 5.22 Geo 6.1 6.2 6.3 6.4 6.5 6.6	L m/s   Enumeration   Strling Number   5.2.1 First Kind   5.2.2 Second Kind   ax+by=gcd   Berlekamp Massey   Charateristic Polynomial   Chinese Remainder   De-Bruijn   DiscreteLog   Extended Euler   Extended FloorSum   Fast Fourier Transform   FloorSum   FWT   Miller Rabin   NTT   Partition Number   Pi Count (Linear Sieve)   Pollard Rho   Polynomial Operations   Quadratic residue   Simplex   2 Simplex Construction   cometry   Basic Geometry   Segment & Line Intersection   2D Convex Hull   3D Convex Hull   2D Farthest Pair   kD Closest Pair (3D ver.)	
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 5.21 5.22 6.3 6.4 6.5 6.5 6.5 6.6 6.7	L m/s   Enumeration   Strling Number   5.2.1 First Kind   5.2.2 Second Kind   ax+by=gcd   Berlekamp Massey   Charateristic Polynomial   Chinese Remainder   De-Bruijn   DiscreteLog   Extended Euler   Extended FloorSum   Fast Fourier Transform   FloorSum   FWT   Miller Rabin   NTT   Partition Number   Pi Count (Linear Sieve)   Pollard Rho   Polynomial Operations   Quadratic residue   Simplex   2 Simplex Construction   Ometry   Basic Geometry   Segment & Line Intersection   2D Convex Hull   3D Convex Hull   3D Convex Hull   3D Farthest Pair   kD Closest Pair (3D ver.)   Simulated Annealing	
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.15 5.16 5.17 5.18 5.20 5.21 5.22 6.3 6.4 6.5 6.6 6.7 6.6 6.7 6.7 6.7 6.7 6.7 6.7 6.7	L	
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	5.1 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.10 6.11 6.12 6.13 6.14 6.14 6.15 6.16 6.17 6.17 6.17 6.17 6.17 6.17 6.17	Times   Enumeration   Strling Number   5.2.1 First Kind   5.2.2 Second Kind   αx+by=gcd   Berlekamp Massey   Charateristic Polynomial   Chinese Remainder   De-Bruijn   DiscreteLog   Extended Euler   ExtendedFloorSum   Fast Fourier Transform   FloorSum   FWT   Miller Rabin   NTT   Partition Number   Pi Count (Linear Sieve)   Pollard Rho   Polynomial Operations   Quadratic residue   Simplex   Simplex   Simplex Construction   Dometry   Basic Geometry   Segment & Line Intersection   2D Convex Hull   3D Convex Hull   3D Convex Hull   2D Farthest Pair   KD Closest Pair (3D ver.)   Simulated Annealing   Half Plane Intersection   Minkowski Sum   Circle Class   Intersection of line and Circle   Intersection of Polygon and Circle   Point & Hulls Tangent   Convex Hulls Tangent	
	5.1 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.14 5.15 5.16 5.17 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.10 6.11 6.12 6.14 6.15 6.14 6.15 6.16 6.17 6.17 6.18 6.19 6.19 6.19 6.19 6.19 6.19 6.19 6.19	Part   Enumeration   Strling Number   5.2.1 First Kind   5.2.2 Second Kind   αx+by=gcd   Berlekamp Massey   Charateristic Polynomial   Chinese Remainder   De-Bruijn   DiscreteLog   Extended Euler   Extended FloorSum   Fast Fourier Transform   FloorSum   FWT   Miller Rabin   NTT   Partition Number   Pi Count (Linear Sieve)   Pollard Rho   Polynomial Operations   Quadratic residue   Simplex   Simplex   Simplex Construction   Desgment & Line Intersection   DD Convex Hull   DD Farthest Pair   KD Closest Pair (3D ver.)   Simulated Annealing   Half Plane Intersection   Minkowski Sum   Circle Class   Intersection of Polygon and Circle   Point & Hulls Tangent   Convex Hulls Tangent   Co	
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```
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8 Misc
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     8.1.2
  8.1.3 Tutte's Matrix
8.1.4 Cayley's Formula
8.1.5 Erdős–Gallai theorem
8.1.6 Havel–Hakimi algorithm
8.1.7 Euler's planar graph formula
  Basic
1.1 vimrc
se is nu ru et tgc sc hls cin cino+=j1 sw=4 sts=4 bs=2
  mouse=a encoding=utf-8 ls=2
syn on
colorscheme desert
filetype indent on
inoremap {<CR> {<CR>}<ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -std=c++17 -
  DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
  Wconversion -fsanitize=address,undefined -g && echo
  success<CR>
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 &&
  echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
```

# 1.2 Debug Macro

```
#ifdef KISEKI
#define safe cerr<<__PRETTY_FUNCTION__\</pre>
   <<" line "<<__LINE__<<" safe\n'
#define debug(a...) qwerty(#a, a)
#define orange(a...) dvorak(#a, a)
using std::cerr;
template <typename ...T>
void qwerty(const char *s, T ...a) {
  cerr << "\e[1;32m(" << s << ") = (";</pre>
   int cnt = sizeof...(T);
   (..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
template <typename Iter>
void dvorak(const char *s, Iter L, Iter R) {
  cerr << "\e[1;32m[ " << s << " ] = [ ";</pre>
  for (int f = 0; L != R; ++L)
    cerr << (f++ ? ", " : "") << *L;
    cerr << " ]\e[0m\n";
#else
#define safe ((void)0)
#define debug(...) ((void)0)
#define orange(...) ((void)0)
#endif
```

# 1.3 Increase Stack

```
const int size = 256 << 20;
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));</pre>
```

### 1.4 Pragma Optimization

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,sse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8000)
:
```

```
1.5 IO Optimization
```

```
static inline int gc() {
  constexpr int B = 1<<20;
  static char buf[B], *p, *q;
  if(p == q &&
    (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
    return EOF;
  return *p++;
}

template < typename T >
  static inline bool gn( T &x ) {
  int c = gc(); T sgn = 1; x = 0;
  while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
  if(c == '-') sgn = -1, c = gc();
  if(c == EOF) return false;
  while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
  return x *= sgn, true;
}</pre>
```

# 2 Data Structure

# 2.1 Dark Magic

```
2.2 Link-Cut Tree
template <typename Val> class LCT {
private:
struct node
 int pa, ch[2];
 bool rev;
 Val v, v_prod, v_rprod;
 node() : pa{0}, ch{0, 0}, rev{false}, v{}, v_prod{},
    v_rprod{} {};
vector<node> nodes;
 set<pair<int, int>> edges;
bool is_root(int u) const {
 const int p = nodes[u].pa;
  return nodes[p].ch[0] != u and nodes[p].ch[1] != u;
bool is_rch(int u) const {
 return (not is_root(u)) and nodes[nodes[u].pa].ch[1]
    == u;
void down(int u) {
 if (auto &cnode = nodes[u]; cnode.rev) {
  if (cnode.ch[0]) set_rev(cnode.ch[0]);
   if (cnode.ch[1]) set_rev(cnode.ch[1]);
   cnode.rev = false;
}
void up(int u) {
 auto &cnode = nodes[u];
 cnode.v_prod =
  nodes[cnode.ch[0]].v_prod * cnode.v * nodes[cnode.ch
    [1]].v_prod;
 cnode.v_rprod =
  nodes[cnode.ch[1]].v_rprod * cnode.v * nodes[cnode.
    ch[0]].v_rprod;
void set_rev(int u) {
 swap(nodes[u].ch[0],\ nodes[u].ch[1]);\\
  swap(nodes[u].v_prod, nodes[u].v_rprod);
 nodes[u].rev ^= 1;
 void rotate(int u) {
 int f = nodes[u].pa, g = nodes[f].pa, l = is_rch(u);
```

```
if (nodes[u].ch[1 ^ 1])
   nodes[nodes[u].ch[1 ^ 1]].pa = f;
  if (not is_root(f))
   nodes[g].ch[is_rch(f)] = u;
  nodes[f].ch[1] = nodes[u].ch[1 ^ 1];
  nodes[u].ch[1 ^ 1] = f
  nodes[u].pa = g, nodes[f].pa = u;
  up(f);
 void splay(int u) {
  vector<int> stk = {u};
  while (not is_root(stk.back()))
   stk.push_back(nodes[stk.back()].pa);
  for (; not stk.empty(); stk.pop_back())
   down(stk.back());
  for(int f=nodes[u].pa;!is_root(u);f=nodes[u].pa){
   if(!is_root(f))rotate(is_rch(u)==is_rch(f)?f:u);
   rotate(u):
  }
  up(u);
 void access(int u) {
  int last = 0;
  for (int last = 0; u; last = u, u = nodes[u].pa) {
   splay(u);
   nodes[u].ch[1] = last;
   up(u);
 int find_root(int u) {
  access(u); splay(u);
  int la = 0:
  for (; u; la = u, u = nodes[u].ch[0]) down(u);
  return la;
 void change_root(int u) {
  access(u); splay(u); set_rev(u);
 void link(int x, int y)
  change_root(y); nodes[y].pa = x;
 void split(int x, int y) {
  change_root(x); access(y); splay(y);
 void cut(int x, int y) {
  split(x, y)
  nodes[y].ch[0] = nodes[x].pa = 0;
  up(y);
public:
 LCT(int n = 0) : nodes(n + 1) {}
 int add(const Val &v = {}) {
 nodes.push_back(v);
  return int(nodes.size()) - 2;
 int add(Val &&v) {
  nodes.emplace_back(move(v));
  return int(nodes.size()) - 2;
 void set_val(int u, const Val &v) {
  splay(++u); nodes[u].v = v; up(u);
 Val query(int x, int y) {
  split(++x, ++y);
  return nodes[y].v_prod;
 bool connected(int u, int v) { return find_root(++u)
    == find_root(++v);
 void add_edge(int u, int v) {
  if (++u > ++v) swap(u, v)
  edges.emplace(u, v); link(u, v);
 void del_edge(int u, int v) {
  auto k = minmax(++u, ++v)
  if (auto it = edges.find(k); it != edges.end()) {
   edges.erase(it); cut(u, v);
  }
};
```

# .3 LiChao Segment Tree

struct L {

```
int m, k, id;
                                                                #undef sz
 L(): id(-1) {}
 L(int a, int b, int c) : m(a), k(b), id(c) {}
                                                               2.5 Linear Basis
 int at(int x) { return m * x + k; }
                                                               template <int BITS>
class LiChao {
                                                               struct LinearBasis {
private:
                                                                array<uint64_t, BITS> basis;
 int n; vector<L> nodes;
                                                                Basis() { basis.fill(0); }
                                                                void add(uint64_t x)
 static int lc(int x) { return 2 * x + 1; }
 static int rc(int x) { return 2 * x + 2;
                                                                 for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
  if (basis[i] == 0) {
 void insert(int 1, int r, int id, L ln) {
  int m = (1 + r) >> 1;
                                                                   basis[i] = x;
  if (nodes[id].id == -1) {
                                                                    return;
   nodes[id] = ln;
   return;
                                                                  x ^= basis[i];
                                                                 }
  bool atLeft = nodes[id].at(1) < ln.at(1);</pre>
  if (nodes[id].at(m) < ln.at(m)) {</pre>
                                                                bool ok(uint64_t x) {
   atLeft ^= 1;
                                                                 for (int i = 0; i < BITS; ++i)</pre>
   swap(nodes[id], ln);
                                                                  if ((x >> i) & 1) x ^= basis[i];
                                                                 return x == 0;
  if (r - 1 == 1) return;
  if (atLeft) insert(1, m, lc(id), ln);
                                                               };
  else insert(m, r, rc(id), ln);
                                                               2.6
                                                                      Binary Search On Segment Tree
 int query(int 1, int r, int id, int x) {
  int ret = 0, m = (1 + r) >> 1;
                                                               // find_first = x -> minimal x s.t. check( [a, x) )
                                                               // find_last = x \rightarrow maximal x s.t. check([x, b))
  if (nodes[id].id != -1)
                                                               template <typename C>
   ret = nodes[id].at(x);
                                                               int find_first(int 1, const C &check) {
  if (r - 1 == 1) return ret;
                                                                if (1 >= n) return n + 1;
  if (x < m) return max(ret, query(1, m, lc(id), x));</pre>
                                                                1 += sz;
  return max(ret, query(m, r, rc(id), x));
                                                                for (int i = height; i > 0; i--)
                                                                 propagate(1 >> i);
                                                                Monoid sum = identity;
public:
 LiChao(int n_{-}) : n(n_{-}), nodes(n * 4) {}
                                                                 while ((1 & 1) == 0) 1 >>= 1;
 void insert(L ln) { insert(0, n, 0, ln); }
                                                                 if (check(f(sum, data[1]))) {
 int query(int x) { return query(0, n, 0, x); }
                                                                  while (1 < sz) {</pre>
                                                                   propagate(1);
                                                                    1 <<= 1;
2.4 Treap
                                                                    auto nxt = f(sum, data[1]);
namespace Treap{
                                                                    if (not check(nxt)) {
 #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                                     sum = nxt;
 struct node{
                                                                     1++;
  int size;
                                                                   }
  uint32_t pri;
                                                                  }
  node *lc, *rc, *pa;
                                                                   return 1 + 1 - sz;
  node():size(0),pri(rand()),lc(0),rc(0),pa(0){}
  void pull() {
                                                                 sum = f(sum, data[1++]);
  size = 1; pa = nullptr;
                                                                } while ((1 & -1) != 1);
   if ( lc ) { size += lc->size; lc->pa = this; }
if ( rc ) { size += rc->size; rc->pa = this; }
                                                                return n + 1;
  }
                                                               template <typename C>
                                                               int find_last(int r, const C &check) {
node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
                                                                if (r <= 0) return -1;
                                                                r += sz;
  if ( L->pri > R->pri ) {
                                                                for (int i = height; i > 0; i--)
  L->rc = merge( L->rc, R ); L->pull();
                                                                 propagate((r - 1) >> i);
   return L;
                                                                Monoid sum = identity;
  } else {
                                                                do {
   R->lc = merge( L, R->lc ); R->pull();
   return R;
                                                                 while (r > 1 \text{ and } (r \& 1)) r >>= 1;
  }
                                                                 if (check(f(data[r], sum))) {
 }
                                                                  while (r < sz) {
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                                   propagate(r);
  if ( not rt ) L = R = nullptr;
                                                                    r = (r << 1) + 1;
  else if( sz( rt->lc ) + 1 <= k ) {
                                                                    auto nxt = f(data[r], sum);
                                                                    if (not check(nxt)) {
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                     sum = nxt;
   L->pull();
                                                                     r--:
  } else {
                                                                   }
   R = rt:
   split_by_size( rt->lc, k, L, R->lc );
                                                                  return r - sz;
   R->pull();
                                                                 }
                                                                sum = f(data[r], sum);
} while ((r & -r) != r);
 } // sz(L) == k
 int getRank(node *o) { // 1-base
int r = sz(o->lc) + 1;
                                                                return -1;
  for (;o->pa != nullptr; o = o->pa)
   if (o->pa->rc == o) r += sz(o->pa->lc) + 1;
                                                                    Graph
  return r;
```

3.1 2-SAT (SCC)

G.clear(); G.resize(n = n\_);

```
class TwoSat{
                                                                 low.assign(n, ecnt = 0);
private:
                                                                dfn.assign(n, 0);
 int n;
                                                               }
 vector<vector<int>> rG,G,sccs;
                                                               void add_edge(int u, int v) {
                                                                G[u].emplace_back(v, ecnt);
 vector<int> ord,idx;
 vector<bool> vis,result;
                                                                G[v].emplace_back(u, ecnt++);
 void dfs(int u){
                                                               void solve() {
  vis[u]=true
  for(int v:G[u])
                                                                bridge.assign(ecnt, false);
                                                                for (int i = 0; i < n; ++i)
   if(!vis[v]) dfs(v);
                                                                 if (not dfn[i]) dfs(i, i);
  ord.push_back(u);
 void rdfs(int u){
                                                               bool is_bridge(int x) { return bridge[x]; }
                                                             } bcc_bridge;
 vis[u]=false;idx[u]=sccs.size()-1;
  sccs.back().push_back(u);
                                                             3.3 BCC Vertex
  for(int v:rG[u])
   if(vis[v])rdfs(v);
                                                             class BCC_AP {
                                                              private:
public:
                                                               int n, ecnt;
 void init(int n_){
                                                               vector<vector<pair<int,int>>> G;
                                                               vector<int> bcc, dfn, low, st;
 G.clear();G.resize(n=n_);
                                                               vector<bool> ap, ins;
void dfs(int u, int f)
  rG.clear();rG.resize(n);
  sccs.clear();ord.clear();
                                                                dfn[u] = low[u] = dfn[f] + 1;
  idx.resize(n);result.resize(n);
                                                                 int ch = 0;
                                                                for (auto [v, t]: G[u]) if (v != f) {
  if (not ins[t]) {
 void add_edge(int u,int v){
 G[u].push_back(v);rG[v].push_back(u);
                                                                  st.push_back(t);
                                                                   ins[t] = true;
 void orr(int x,int y){
  if ((x^y)==1)return
  add_edge(x^1,y); add_edge(y^1,x);
                                                                  if (dfn[v]) {
                                                                  low[u] = min(low[u], dfn[v]);
 bool solve(){
                                                                   continue:
  vis.clear();vis.resize(n);
                                                                  } ++ch; dfs(v, u);
  for(int i=0;i<n;++i)</pre>
                                                                  low[u] = min(low[u], low[v]);
   if(not vis[i])dfs(i);
                                                                  if (low[v] >= dfn[u]) {
                                                                  ap[u] = true;
  reverse(ord.begin(),ord.end());
  for (int u:ord){
                                                                   while (true) {
   if(!vis[u])continue;
                                                                    int eid = st.back(); st.pop_back();
   sccs.push_back(vector<int>());
                                                                    bcc[eid] = ecnt;
   rdfs(u);
                                                                    if (eid == t) break;
  for(int i=0;i<n;i+=2)</pre>
                                                                  ecnt++;
   if(idx[i]==idx[i+1])
                                                                 }
    return false;
  vector<bool> c(sccs.size());
                                                                 if (ch == 1 and u == f) ap[u] = false;
  for(size_t i=0;i<sccs.size();++i){</pre>
                                                              public:
   for(auto sij : sccs[i]){
    result[sij]=c[i];
                                                               void init(int n_) {
    c[idx[sij^1]]=!c[i];
                                                                G.clear(); G.resize(n = n_);
   }
                                                                 ecnt = 0; ap.assign(n, false);
                                                                low.assign(n, 0); dfn.assign(n, 0);
  return true;
                                                               void add_edge(int u, int v) {
                                                                G[u].emplace_back(v, ecnt);
G[v].emplace_back(u, ecnt++);
 bool get(int x){return result[x];}
 int get_id(int x){return idx[x];}
 int count(){return sccs.size();}
} sat2;
                                                               void solve() {
                                                                ins.assign(ecnt, false);
      BCC Edge
3.2
                                                                bcc.resize(ecnt); ecnt = 0;
                                                                for (int i = 0; i < n; ++i)
if (not dfn[i]) dfs(i, i);</pre>
class BCC_Bridge {
 private:
  int n, ecnt;
                                                               int get_id(int x) { return bcc[x]; }
  vector<vector<pair<int,int>>> G;
  vector<int> dfn, low;
                                                               int count() { return ecnt;
  vector<bool> bridge;
                                                               bool is_ap(int x) { return ap[x]; }
  void dfs(int u, int f)
                                                             } bcc_ap;
   dfn[u] = low[u] = dfn[f] + 1;
                                                             3.4 Centroid Decomposition
   for (auto [v, t]: G[u]) {
    if (v == f) continue;
                                                             struct Centroid {
    if (dfn[v]) {
                                                              vector<vector<int64_t>> Dist;
                                                              vector<int> Parent, Depth;
     low[u] = min(low[u], dfn[v]);
     continue;
                                                              vector<int64_t> Sub, Sub2;
                                                              vector<int> Sz, Sz2;
    dfs(v, u);
                                                              Centroid(vector<vector<pair<int, int>>> g) {
    low[u] = min(low[u], low[v]);
                                                               int N = g.size();
                                                               vector<bool> Vis(N);
    if (low[v] > dfn[u]) bridge[t] = true;
                                                               vector<int> sz(N), mx(N);
                                                               vector<int> Path;
                                                               Dist.resize(N)
 public:
  void init(int n_) {
                                                               Parent.resize(N);
```

Depth.resize(N);

in[e.v] = e.w;

```
auto DfsSz = [&](auto dfs, int x) -> void {
                                                                   prv[e.v] = e.u;
   Vis[x] = true; sz[x] = 1; mx[x] = 0;
   for (auto [u, w] : g[x]) {
                                                                 in[root] = 0;
    if (Vis[u]) continue;
                                                                 prv[root] = -1;
    dfs(dfs, u)
                                                                 for (int i = 0; i < n; i++)
                                                                  if (in[i] == -inf)
    sz[x] += sz[u];
    mx[x] = max(mx[x], sz[u]);
                                                                   return -inf;
                                                                  // find cycle
   Path.push_back(x);
                                                                  int tot = 0;
                                                                 vector<int> id(n, -1), vis(n, -1);
for (int i = 0; i < n; i++) {</pre>
  }:
  auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
                                                                  ans += in[i];
   Dist[x].push_back(D);Vis[x] = true;
                                                                  for (int x = i; x != -1 && id[x] == -1; x = prv[x])
   for (auto [u, w] : g[x]) {
   if (Vis[u]) continue;
                                                                   if (vis[x] == i) {
    dfs(dfs, u, D + w);
                                                                     for (int y = prv[x]; y != x; y = prv[y])
                                                                      id[y] = tot;
  };
                                                                     id[x] = tot++;
  auto Dfs = [&]
                                                                    break;
   (auto dfs, int x, int D = 0, int p = -1)->void {
   Path.clear(); DfsSz(DfsSz, x);
                                                                   vis[x] = i;
   int M = Path.size();
                                                                  }
   int C = -1;
   for (int u : Path) {
                                                                 if (!tot)
                                                                  return ans;
    if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
    Vis[u] = false;
                                                                 for (int i = 0; i < n; i++)</pre>
                                                                  if (id[i] == -1)
   DfsDist(DfsDist, C);
                                                                   id[i] = tot++;
   for (int u : Path) Vis[u] = false;
                                                                  // shrink
   Parent[C] = p; Vis[C] = true;
                                                                 for (auto &e : E) {
   Depth[C] = D;
                                                                  if (id[e.u] != id[e.v])
   for (auto [u, w] : g[C]) {
                                                                   e.w -= in[e.v];
    if (Vis[u]) continue
                                                                  e.u = id[e.u], e.v = id[e.v];
    dfs(dfs, u, D + 1, C);
                                                                 n = tot;
                                                                 root = id[root];
  Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
  Sz.resize(N); Sz2.resize(N);
                                                                assert(false);
void Mark(int v) {
                                                              } DMST:
  int x = v, z = -1
                                                              3.6 Dominator Tree
 for (int i = Depth[v]; i >= 0; --i) {
Sub[x] += Dist[v][i]; Sz[x]++;
                                                              namespace dominator {
  if (z != -1) {
                                                              vector<int> g[maxn], r[maxn], rdom[maxn];
                                                              int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
    Sub2[z] += Dist[v][i];
    Sz2[z]++;
                                                              void init(int n) {
   z = x; x = Parent[x];
                                                               // vertices are numbered from 0 to n-1
  }
                                                               fill(dfn, dfn + n, -1);fill(rev, rev + n, -1);
                                                               fill(fa, fa + n, -1); fill(val, val + n, -1);
int64_t Query(int v) {
                                                               fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
                                                               fill(dom, dom + n, -1); tk = 0;
for (int i = 0; i < n; ++i) {
 int64_t res = 0;
 int x = v, z = -1;
 for (int i = Depth[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
                                                                g[i].clear(); r[i].clear(); rdom[i].clear();
  if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
  z = x; x = Parent[x];
                                                              void add_edge(int x, int y) { g[x].push_back(y); }
                                                              void dfs(int x) {
                                                               rev[dfn[x] = tk] = x;
  return res;
                                                               fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
                                                               for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
};
3.5 Directed Minimum Spanning Tree
                                                                r[dfn[u]].push_back(dfn[x]);
struct DirectedMST { // find maximum
                                                               }
struct Edge {
                                                              void merge(int x, int y) { fa[x] = y; }
 int u, v;
                                                              int find(int x, int c = 0) {
  int w;
                                                               if (fa[x] == x) return c ? -1 : x;
 Edge(int u, int v, int w) : u(u), v(v), w(w) {}
                                                               int p = find(fa[x], 1);
                                                               if (p == -1) return c ? fa[x] : val[x];
vector<Edge> Edges;
                                                               if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
void clear() { Edges.clear(); }
void addEdge(int a, int b, int w) { Edges.emplace_back
                                                               fa[x] = p
                                                               return c ? p : val[x];
    (a, b, w); }
int solve(int root, int n) {
  vector<Edge> E = Edges;
                                                              vector<int> build(int s, int n) {
                                                              // return the father of each node in the dominator tree
  int ans = 0:
 while (true) {
                                                              // p[i] = -2 if i is unreachable from s
   // find best in edge
                                                               dfs(s);
   vector<int> in(n, -inf), prv(n, -1);
                                                               for (int i = tk - 1; i >= 0; --i) {
                                                                for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
                                                                if (i) rdom[sdom[i]].push_back(i);
```

for (int &u : rdom[i]) {

chain[u] = chain[v];

```
int p = find(u);
   if (sdom[p] == i) dom[u] = i;
                                                                    if (chain[u] == 0) chain[u] = ++chains;
   else dom[u] = p;
                                                                   void dfschain(int u, int f) {
                                                                    tl[u] = timer++
 if (i) merge(i, rp[i]);
                                                                    if (head[chain[u]] == -1)
vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)</pre>
                                                                     head[chain[u]] = u;
                                                                    for (int v : G[u])
  if (v != f and chain[v] == chain[u])
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
                                                                      dfschain(v, u);
                                                                    for (int v : G[u])
 return p;
                                                                     if (v != f and chain[v] != chain[u])
                                                                      dfschain(v, u);
3.7 Edge Coloring
                                                                    tr[u] = timer;
// max(d_u) + 1 edge coloring, time: O(NM)
                                                                 public:
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
                                                                   LBD(int n) : timer(0), chains(0), G(n), tl(n), tr(n),
void clear(int N) {
                                                                      chain(n), \ head(n, \ -1), \ dep(n), \ pa(n) \ \{\}
for (int i = 0; i <= N; i++)
 for (int j = 0; j <= N; j++)
C[i][j] = G[i][j] = 0;
                                                                   void add_edge(int u, int v) {
                                                                    G[u].push_back(v); G[v].push_back(u);
void solve(vector<pair<int, int>> &E, int N) {
                                                                   void decompose() { predfs(0, 0); dfschain(0, 0); }
int X[kN] = {}, a;
auto update = [&](int u) {
                                                                   PII get_subtree(int u) { return {tl[u], tr[u]}; }
                                                                   vector<PII> get_path(int u, int v) {
                                                                    vector<PII> res;
while (chain[u] != chain[v]) {
 for (X[u] = 1; C[u][X[u]]; X[u]++);
auto color = [&](int u, int v, int c) {
                                                                     if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
 int p = G[u][v];
                                                                      swap(u, v)
                                                                     int s = head[chain[u]];
res.emplace_back(tl[s], tl[u] + 1);
 G[u][v] = G[v][u] = c;
 C[u][c] = v, C[v][c] = u;
 C[u][p] = C[v][p] = 0;
                                                                     u = pa[s];
 if (p) X[u] = X[v] = p;
  else update(u), update(v);
                                                                    if (dep[u] < dep[v]) swap(u, v);</pre>
                                                                    res.emplace_back(tl[v], tl[u] + 1);
  return p;
                                                                    return res;
auto flip = [&](int u, int c1, int c2) {
 int p = C[u][c1];
                                                                 };
 swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
if (!C[u][c1]) X[u] = c1;
                                                                        Manhattan Minimum Spanning Tree
                                                                 typedef Point<int> P:
  if (!C[u][c2]) X[u] = c2;
                                                                 vector<array<int, 3>> manhattanMST(vector<P> ps) {
                                                                   vi id(sz(ps));
 return p;
                                                                   iota(all(id), 0);
for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {
  auto [u, v] = E[t];</pre>
                                                                   vector<array<int, 3>> edges;
                                                                   rep(k, 0, 4) {
  sort(all(id),
                                                                                   [&](int i, int j) {
  int v0 = v, c = X[u], c0 = c, d;
                                                                     return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
 vector<pair<int, int>> L; int vst[kN] = {};
                                                                    });
  while (!G[u][v0]) {
                                                                    map<int, int> sweep;
   L.emplace_back(v, d = X[v]);
if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
                                                                    for (int i : id) {
                                                                     for (auto it = sweep.lower_bound(-ps[i].y);
     c = color(u, L[a].first, c);
                                                                        it != sweep.end(); sweep.erase(it++)) {
                                                                      int j = it->second
   else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
     color(u, L[a].first, L[a].second);
                                                                      P d = ps[i] - ps[j];
   else if (vst[d]) break
                                                                      if (d.y > d.x) break;
   else vst[d] = 1, v = C[u][d];
                                                                      edges.push_back({d.y + d.x, i, j});
 if (!G[u][v0]) {
  for (; v; v = flip(v, c, d), swap(c, d));
                                                                     sweep[-ps[i].y] = i;
   if (C[u][c0]) { a = int(L.size()) - 1;
                                                                    for (P &p : ps)
    while (--a >= 0 && L[a].second != c);
                                                                     if (k \& 1) p.x = -p.x;
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
                                                                     else swap(p.x, p.y);
   } else t--;
                                                                   return edges; // [{w, i, j}, ...]
                                                                 }
                                                                 3.10 MaxClique
3.8 Lowbit Decomposition
                                                                 // contain a self loop u to u, than u won't in clique
                                                                 template < size_t MAXN >
class LBD {
int timer, chains;
                                                                 class MaxClique{
vector<vector<int>> G;
                                                                 private:
vector<int> t1, tr, chain, head, dep, pa;
                                                                   using bits = bitset< MAXN >;
 // chains : number of chain
                                                                   bits popped, G[ MAXN ], ans;
// tl, tr[u] : subtree interval in the seq. of u
                                                                   size_t deg[ MAXN ], deo[ MAXN ], n;
 // head[i] : head of the chain i
                                                                   void sort_by_degree() {
 // chian[u] : chain id of the chain u is on
                                                                    popped.reset();
void predfs(int u, int f) {
  dep[u] = dep[pa[u] = f] + 1;
                                                                    for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                      deg[ i ] = G[ i ].count();
                                                                    for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
  for (int v : G[u]) if (v != f) {
                                                                      size_t mi = MAXN, id = 0;
   predfs(v, u);
                                                                      for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )</pre>
   if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
```

p < kN; p = int(cs[k].\_Find\_next(p))) {

```
mi = deg[ id = j ];
                                                                    r[t] = p; c[t++] = k;
    popped[ deo[ i ] = id ] = 1;
    for( size_t u = G[ i ]._Find_first() ;
     u < n ; u = G[ i ]._Find_next( u ) )</pre>
      -- deg[ u ];
                                                                 void dfs(vector<int> &r, vector<int> &c, int 1,
  }
                                                                  bitset<kN> mask) {
                                                                  while (!r.empty()) {
 void BK( bits R, bits P, bits X ) {
                                                                   int p = r.back(); r.pop_back();
  if (R.count()+P.count() <= ans.count()) return;</pre>
                                                                   mask[p] = 0;
  if ( not P.count() and not X.count() ) {
                                                                   if (q + c.back() <= ans) return;</pre>
  if ( R.count() > ans.count() ) ans = R;
                                                                   cur[q++] = p;
                                                                   vector<int> nr, nc;
   return;
                                                                   bitset<kN> nmask = mask & a[p];
  }
  /* greedily chosse max degree as pivot
                                                                   for (int i : r)
  bits cur = P | X; size_t pivot = 0, sz = 0;
                                                                    if (a[p][i]) nr.push_back(i);
  for ( size_t u = cur._Find_first() ;
                                                                   if (!nr.empty()) {
   u < n ; u = cur._Find_next( u )</pre>
                                                                    if (1 < 4) {
   if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
                                                                     for (int i : nr)
  cur = P & ( ~G[ pivot ] );
                                                                       d[i] = int((a[i] & nmask).count());
  */ // or simply choose first
                                                                      sort(nr.begin(), nr.end(),
  bits cur = P & (~G[ ( P | X )._Find_first() ]);
                                                                       [&](int x, int y)
  for ( size_t u = cur._Find_first()
                                                                        return d[x] > d[y];
   u < n ; u = cur._Find_next( u ) ) {
if ( R[ u ] ) continue;</pre>
                                                                       });
                                                                   csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
   R[u] = 1;
   BK(R, P & G[u], X & G[u]);
R[u] = P[u] = 0, X[u] = 1;
                                                                    ans = q; copy(cur, cur + q, sol);
                                                                   c.pop_back(); q--;
public:
                                                                  }
 void init( size_t n_ ) {
                                                                 int solve(bitset<kN> mask) { // vertex mask
  n = n_{-};
                                                                  vector<int> r, c;
for (int i = 0; i < n; i++)</pre>
  for ( size_t i = 0 ; i < n ; ++ i )
   G[ i ].reset();
  ans.reset();
                                                                   if (mask[i]) r.push_back(i);
                                                                  for (int i = 0; i < n; i++)</pre>
                                                                   d[i] = int((a[i] & mask).count());
 void add_edges( int u, bits S ) { G[ u ] = S; }
 void add_edge( int u, int v ) {
                                                                  sort(r.begin(), r.end(),
 G[u][v] = G[v][u] = 1;
                                                                   [&](int i, int j) { return d[i] > d[j]; });
                                                                  csort(r, c);
 int solve() {
                                                                  dfs(r, c, 1, mask);
                                                                  return ans; // sol[0 ~ ans-1]
  sort_by_degree(); // or simply iota( deo... )
  for ( size_t i = 0 ; i < n ; ++ i )
   deg[ i ] = G[ i ].count();
                                                               } graph;
  bits pob, nob = 0; pob.set();
                                                                3.12 Minimum Mean Cycle
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
  for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
                                                                /* minimum mean cycle O(VE) */
   size_t v = deo[ i ];
                                                                struct MMC{
   bits tmp; tmp[v] = 1;
                                                                #define FZ(n) memset((n),0,sizeof(n))
   BK( tmp, pob & G[ v ], nob & G[ v ] );
                                                                #define E 101010
   pob[ v ] = 0, nob[ v ] = 1;
                                                                #define V 1021
                                                                #define inf 1e9
                                                                 struct Edge { int v,u; double c; };
  return static_cast< int >( ans.count() );
                                                                 int n, m, prv[V][V], prve[V][V], vst[V];
};
                                                                 Edge e[E];
                                                                 vector<int> edgeID, cycle, rho;
3.11 MaxCliqueDyn
                                                                 double d[V][V];
                                                                 void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
constexpr int kN = 150;
struct MaxClique { // Maximum Clique
bitset<kN> a[kN], cs[kN];
                                                                 void add_edge( int vi , int ui , double ci )
                                                                 { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
 int ans, sol[kN], q, cur[kN], d[kN], n;
 void init(int _n) {
 n = n, and q = 0;
                                                                  for(int i=0; i<n; i++) d[0][i]=0;</pre>
 for (int i = 0; i < n; i++) a[i].reset();</pre>
                                                                  for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);
  for(int j=0; j<m; j++) {</pre>
 void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
 void csort(vector<int> &r, vector<int> &c)
                                                                    int v = e[j].v, u = e[j].u;
                                                                    if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
  d[i+1][u] = d[i][v]+e[j].c;
  int mx = 1, km = max(ans - q + 1, 1), t = 0,
m = int(r.size());
  cs[1].reset(); cs[2].reset();
                                                                     prv[i+1][u] = v;
  for (int i = 0; i < m; i++) {
  int p = r[i], k = 1;</pre>
                                                                     prve[i+1][u] = j;
   while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
                                                                  }
   cs[k][p] = 1;
   if (k < km) r[t++] = p;
                                                                 double solve(){
                                                                  // returns inf if no cycle, mmc otherwise
  c.resize(m);
                                                                  double mmc=inf;
  if(t) c[t-1] = 0;
                                                                  int st = -1;
  for (int k = km; k <= mx; k++) {</pre>
                                                                  bellman_ford();
   for (int p = int(cs[k]._Find_first());
                                                                  for(int i=0; i<n; i++) {</pre>
```

double avg=-inf;

```
for(int k=0; k<n; k++) {</pre>
    if(d[n][i]<inf-eps)</pre>
     avg=max(avg,(d[n][i]-d[k][i])/(n-k));
    else avg=max(avg,inf);
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
 FZ(vst);edgeID.clear();cycle.clear();rho.clear();
  for (int i=n; !vst[st]; st=prv[i--][st]) {
  vst[st]++
   edgeID.PB(prve[i][st]);
   rho.PB(st);
 while (vst[st] != 2) {
  int v = rho.back(); rho.pop_back();
   cycle.PB(v);
   vst[v]++;
 reverse(ALL(edgeID));
 edgeID.resize(SZ(cycle));
  return mmc:
} mmc;
3.13
      Minimum Steiner Tree
```

```
// Minimum Steiner Tree
// 0(V 3^T + V^2 2^T)
struct SteinerTree {
#define V 33
#define T 8
#define INF 1023456789
int n, dst[V][V], dp[1 << T][V], tdst[V];</pre>
void init(int _n) {
 for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++)
  dst[i][j] = INF * (i != j);</pre>
  }
void add_edge(int ui, int vi, int wi) {
  dst[ui][vi] = min(dst[ui][vi], wi);
  dst[vi][ui] = min(dst[vi][ui], wi);
void shortest_path() {
  for (int k = 0; k < n; k++)
   for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < n; j++)
dst[i][j] = min(dst[i][j], dst[i][k] + dst[k][j]);</pre>
 int solve(const vector<int> &ter) {
  int t = (int)ter.size();
  for (int i = 1; i < (1 << t); i++)
   fill_n(dp[i], n, INF);
  fill_n(dp[0], n, 0);
  for (int msk = 1; msk < (1 << t); msk++) {
   if (msk == (msk & (-msk))) {
    int who = __lg(msk);
for (int i = 0; i < n; i++)</pre>
     dp[msk][i] = dst[ter[who]][i];
    continue;
   for (int i = 0; i < n; i++)
    for (int submsk = (msk - 1) & msk; submsk; submsk =
      (submsk - 1) & msk)
     dp[msk][i] = min(dp[msk][i], dp[submsk][i] + dp[
    msk ^ submsk][i]);
   for (int i = 0; i < n; i++) {</pre>
    tdst[i] = INF;
    for (int j = 0; j < n; j++)</pre>
     tdst[i] = min(tdst[i], dp[msk][j] + dst[j][i]);
   copy_n(tdst, n, dp[msk]);
  }
  int ans = INF;
  for (int i = 0; i < n; i++)</pre>
   ans = min(ans, dp[(1 << t) - 1][i]);
  return ans;
} solver;
```

# 3.14 Mo's Algorithm on Tree

```
push u
 iterate subtree
 push u
Let P = LCA(u, v) with St(u) <= St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
3.15 Virtual Tree
vector<pair<int, int>> build(vector<int> vs, int r) {
 vector<pair<int, int>> res;
sort(vs.begin(), vs.end(), [](int i, int j) {
  return dfn[i] < dfn[j]; });
 vector<int> s = {r};
 for (int v : vs) if (v != r) {
  if (int o = lca(v, s.back()); o != s.back()) {
   while (s.size() >= 2) {
  if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
    res.emplace_back(s[s.size() - 2], s.back());
    s.pop_back();
   if (s.back() != o) {
    res.emplace_back(s.back(), o);
    s.back() = o;
  s.push_back(v);
 for (size_t i = 1; i < s.size(); ++i)</pre>
  res.emplace_back(s[i - 1], s[i]);
 return res;
4
     Matching & Flow
```

# Bipartite Matching

dfs u:

```
struct BipartiteMatching {
 vector<int> X[N];
 int fX[N], fY[N], n;
 bitset<N> vis;
 bool dfs(int x)
  for (auto i:X[x]) {
   if (vis[i]) continue;
   vis[i] = true;
   if (fY[i]==-1 || dfs(fY[i])){
    fY[fX[x] = i] = x;
    return true;
  return false;
 void init(int n_, int m) {
  vis.reset();
  fill(X, X + (n = n_{-}), vector<int>());
  memset(fX, -1, sizeof(int) * n);
memset(fY, -1, sizeof(int) * m);
 void add_edge(int x, int y){
  X[x].push_back(y); }
 int solve() { // return how many pair matched
  int cnt = 0;
  for(int i=0;i<n;i++) {</pre>
   vis.reset()
   cnt += dfs(i);
  }
  return cnt;
};
4.2 Dijkstra Cost Flow
```

```
// kN = #(vertices)
// MCMF.{Init, AddEdge, MincostMaxflow}
// MincostMaxflow(source, sink, flow_limit, &cost)
// => flow
using Pii = pair<int, int>;
constexpr int kInf = 0x3f3f3f3f, kN = 500;
struct Edge {
int to, rev, cost, flow;
struct MCMF { // 0-based
int n{}, m{}, s{}, t{};
```

```
vector<Edge> graph[kN];
 // Larger range for relabeling
int64_t dis[kN] = {}, h[kN] = {};
 int p[kN] = {};
void Init(int nn) {
 n = nn;
  for (int i = 0; i < n; i++) graph[i].clear();</pre>
 void AddEdge(int u, int v, int f, int c) {
 graph[u].push_back({v,
   static_cast<int>(graph[v].size()), c, f});
  graph[v].push_back(
   {u, static_cast<int>(graph[u].size()) - 1,
    -c, 0});
bool Dijkstra(int &max_flow, int64_t &cost) {
 priority_queue<Pii, vector<Pii>, greater<>> pq;
  fill_n(dis, n, kInf);
  dis[s] = 0;
  pq.emplace(0, s)
  while (!pq.empty()) {
   auto u = pq.top();
   pq.pop();
   int v = u.second;
   if (dis[v] < u.first) continue;</pre>
   for (auto &e : graph[v]) {
    auto new_dis =
     dis[v] + e.cost + h[v] - h[e.to];
    if (e.flow > 0 && dis[e.to] > new_dis) {
     dis[e.to] = new_dis;
     p[e.to] = e.rev;
     pq.emplace(dis[e.to], e.to);
  if (dis[t] == kInf) return false;
  for (int i = 0; i < n; i++) h[i] += dis[i];
  int d = max_flow;
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   d = min(d, graph[e.to][e.rev].flow);
 max_flow -= d;
  cost += int64_t(d) * h[t];
  for (int u = t; u != s
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   e.flow += d;
  graph[e.to][e.rev].flow -= d;
  }
  return true;
int MincostMaxflow(
  int ss, int tt, int max_flow, int64_t &cost) {
 this->s = ss, this->t = tt;
 cost = 0;
 fill_n(h, n, 0);
 auto orig_max_flow = max_flow;
  while (Dijkstra(max_flow, cost) && max_flow) {}
  return orig_max_flow - max_flow;
};
4.3 Dinic
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
    int to, rev;
    Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
      int u = bfs.front(); bfs.pop();
```

for (auto e: G[u]) {

```
if (e.cap <= 0 or lv[e.to]!=-1) continue;
         bfs.push(e.to); lv[e.to] = lv[u] + 1;
       }
     return lv[ed] != -1;
  Cap DFS(int u, Cap f){
     if (u == ed) return f;
     Cap ret = 0;
     for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
       auto &e = G[u][i];
       if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
       Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
       if (f == 0) return ret;
     if (ret == 0) lv[u] = -1;
     return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
     G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
     st = st_, ed = ed_; Cap ret = 0;
     while (BFS()) {
       idx.assign(n, 0);
       Cap f = DFS(st, numeric_limits<Cap>::max());
       ret += f;
       if (f == 0) break;
     return ret;
};
```

#### 4.4 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source  ${\cal S}$  and sink  ${\cal T}$ .
  - 2. For each edge (x,y,l,u), connect  $x \to y$  with capacity u-l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v\to T$  with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, \, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
  - 1. Redirect every edge:  $y \to x$  if  $(x, y) \in M$ ,  $x \to y$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source  ${\cal S}$  and sink  ${\cal T}$
  - 2. For each edge (x,y,c), connect  $x\to y$  with (cost,cap)=(c,1) if c>0, otherwise connect  $y\to x$  with (cost,cap)=(-c,1)
  - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v)>0 , connect  $S\to v$  with (cost,cap)=(0,d(v))
  - 5. For each vertex v with d(v) < 0, connect  $v \to T$  with (cost, cap) = (0, -d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let  ${\cal K}$  be the sum of all weights
  - 3. Connect source  $s \to v$ ,  $v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity
  - 5. For  $v\in G$  , connect it with sink  $v\to t$  with capacity  $K+2T-(\sum_{e\in E(v)}w(e))-2w(v)$
  - 6. T is a valid answer if the maximum flow f < K |V|
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight w(u,v).

- 2. Connect  $v \ \rightarrow \ v'$  with weight  $2\mu(v)\text{, where }\mu(v)$  is the cost of the cheapest edge incident to  $\bar{v}$ .
- 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
  - 1. If  $p_v>0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$
  - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
  - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x, t) with capacity  $c_x$  and create edge (s, y) with capacity  $c_n$
- 2. Create edge (x,y) with capacity  $c_{xy}$ . 3. Create edge (x,y) and edge (x',y') with capacity  $c_{xyx'y'}$ .

# General Graph Matching

```
namespace matching {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
 g[u].push_back(v);
 g[v].push_back(u);
int Find(int u) {
return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
static int tk = 0; tk++;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
void Blossom(int x, int y, int 1) {
 while (Find(x) != 1) {
  pre[x] = y, y = match[x];
  if (s[y] == 1) q.push(y), s[y] = 0;
if (fa[x] == x) fa[x] = 1;
  if (fa[y] == y) fa[y] = 1;
  x = pre[y];
bool Bfs(int r, int n) {
  for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;</pre>
 while (!q.empty()) q.pop();
 q.push(r);
 s[r] = 0;
 while (!q.empty()) {
  int x = q.front(); q.pop();
  for (int u : g[x]) {
   if (s[u] == -1) {
    pre[u] = x, s[u] = 1;
    if (match[u] == n) {
     for (int a = u, b = x, last; b != n; a = last, b =
     pre[a])
      last = match[b], match[b] = a, match[a] = b;
     return true;
    q.push(match[u]);
    s[match[u]] = 0;
   } else if (!s[u] && Find(u) != Find(x)) {
    int 1 = LCA(u, x, n);
Blossom(x, u, 1);
    Blossom(u, x, 1);
```

```
return false;
int Solve(int n) {
 int res = 0;
 for (int x = 0; x < n; ++x) {
  if (match[x] == n) res += Bfs(x, n);
 return res:
4.6 Global Min-Cut
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
pair<int, int> phase(int n) {
memset(v, false, sizeof(v));
 memset(g, 0, sizeof(g));
 int s = -1, t = -1;
 while (true) {
  int c = -1;
  for (int i = 0; i < n; ++i) {
   if (del[i] || v[i]) continue;
   if (c == -1 \mid | g[i] > g[c]) c = i;
  if (c == -1) break;
  v[s = t, t = c] = true;
  for (int i = 0; i < n; ++i) {
   if (del[i] || v[i]) continue;
   g[i] += w[c][i];
 return make_pair(s, t);
int mincut(int n) {
 int cut = 1e9;
 memset(del, false, sizeof(del));
 for (int i = 0; i < n - 1; ++i) {
  int s, t; tie(s, t) = phase(n);
  del[t] = true; cut = min(cut, g[t]);
  for (int j = 0; j < n; ++j) {
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
 return cut;
4.7 GomoryHu Tree
int g[maxn];
vector<edge> GomoryHu(int n){
 vector<edge> rt;
 for(int i=1;i<=n;++i)g[i]=1;</pre>
 for(int i=2;i<=n;++i){</pre>
  int t=g[i];
  flow.reset(); // clear flows on all edge
  rt.push_back({i,t,flow(i,t)});
  flow.walk(i); // bfs points that connected to i (use
    edges not fully flow)
  for(int j=i+1;j<=n;++j){</pre>
   if(g[j]==t && flow.connect(j))g[j]=i; // check if i
    can reach i
 return rt:
4.8 Kuhn Munkres
class KM {
private:
 static constexpr 11d INF = 1LL << 60;</pre>
```

vector<lld> hl,hr,slk; vector<int> fl,fr,pre,qu;

vector<vector<lld>> w;

if (v1[x] = true, f1[x] != -1) return vr[qu[qr++] = f1[x]] = true; while (x != -1) swap(x, fr[fl[x] = pre[x]]);

vector<bool> v1,vr; int n, ql, qr;

bool check(int x) {

```
return false;
                                                                vector<pair<int, int>> cyc;
                                                                while (!mark[rt]) {
 void bfs(int s) {
                                                                 cyc.emplace_back(pv[rt], ed[rt]);
  fill(slk.begin(), slk.end(), INF);
fill(vl.begin(), vl.end(), false);
                                                                 mark[rt] = true;
                                                                 rt = pv[rt];
  fill(vr.begin(), vr.end(), false);
  ql = qr = 0;
                                                                reverse(cyc.begin(), cyc.end());
  vr[qu[qr++] = s] = true;
                                                                int cap = kInf;
  while (true) {
                                                                for (auto &i : cyc)
                                                                 auto &e = g[i.first][i.second];
   11d d;
                                                                 cap = min(cap, e.cap);
   while (ql < qr) {</pre>
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
     if(!v1[x]&&slk[x]>=(d=h1[x]+hr[y]-w[x][y])){
                                                                for (auto &i : cyc)
      if (pre[x] = y, d) slk[x] = d;
                                                                 auto &e = g[i.first][i.second];
      else if (!check(x)) return;
                                                                 e.cap -= cap;
                                                                 g[e.to][e.rev].cap += cap;
                                                                 ans += e.cost * cap;
   d = INF;
                                                               }
   for (int x = 0; x < n; ++x)
                                                               return ans;
    if (!v1[x] \&\& d > s1k[x]) d = s1k[x];
   for (int x = 0; x < n; ++x) {
                                                                    Minimum Cost Maximum Flow
                                                              4.10
    if (v1[x]) h1[x] += d;
    else slk[x] -= d;
                                                             class MiniCostMaxiFlow{
                                                               using Cap = int; using Wei = int64_t;
    if (vr[x]) hr[x] -= d;
                                                               using PCW = pair<Cap,Wei>
   for (int x = 0; x < n; ++x)
                                                               static constexpr Cap INF_CAP = 1 << 30;</pre>
    if (!v1[x] && !slk[x] && !check(x)) return;
                                                               static constexpr Wei INF_WEI = 1LL<<60;</pre>
                                                             private:
                                                               struct Edge{
public:
                                                                int to, back;
 void init( int n_ ) {
                                                                Cap cap; Wei wei;
  qu.resize(n = n_);
fl.assign(n, -1); fr.assign(n, -1);
                                                                Edge() {}
                                                                Edge(int a,int b, Cap c, Wei d):
 hr.assign(n, 0); hl.resize(n);
                                                                 to(a),back(b),cap(c),wei(d) {}
  w.assign(n, vector<lld>(n));
  slk.resize(n); pre.resize(n);
                                                               int ori, edd;
  vl.resize(n); vr.resize(n);
                                                               vector<vector<Edge>> G;
                                                               vector<int> fa, wh;
 void set_edge( int u, int v, lld x ) {w[u][v] = x;}
                                                               vector<bool> inq;
 11d solve() {
                                                               vector<Wei> dis;
 for (int i = 0; i < n; ++i)
                                                               PCW SPFA(){
  hl[i] = *max_element(w[i].begin(), w[i].end());
                                                                fill(inq.begin(),inq.end(),false);
  for (int i = 0; i < n; ++i) bfs(i);</pre>
                                                                fill(dis.begin(), dis.end(), INF_WEI);
  11d res = 0;
                                                                queue<int> qq; qq.push(ori);
  for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
                                                                dis[ori] = 0;
  return res;
                                                                while(not qq.empty()){
 }
                                                                 int u=qq.front();qq.pop();
} km;
                                                                 inq[u] = false;
                                                                 for(int i=0;i<SZ(G[u]);++i){</pre>
     Minimum Cost Circulation
                                                                  Edge e=G[u][i];
struct Edge { int to, cap, rev, cost; };
                                                                  int v=e.to; Wei d=e.wei;
vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
                                                                  if(e.cap <= 0 | |dis[v] <= dis[u] + d)
                                                                   continue
bool mark[kN];
                                                                  dis[v] = dis[u] + d;
int NegativeCycle(int n) {
                                                                  fa[v] = u, wh[v] = i;
memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
                                                                  if (inq[v]) continue;
                                                                  qq.push(v);
 int upd = -1;
                                                                  inq[v] = true;
 for (int i = 0; i <= n; ++i)
  for (int j = 0; j < n; ++j) {
                                                                if(dis[edd]==INF_WEI) return {-1, -1};
   int idx = 0;
   for (auto &e : g[j]) {
                                                                Cap mw=INF_CAP;
    if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
                                                                for(int i=edd;i!=ori;i=fa[i])
     dist[e.to] = dist[j] + e.cost;
                                                                 mw=min(mw,G[fa[i]][wh[i]].cap);
     pv[e.to] = j, ed[e.to] = idx;
                                                                for (int i=edd;i!=ori;i=fa[i]){
     if (i == n) {
                                                                 auto &eg=G[fa[i]][wh[i]];
                                                                 eg.cap
      while(!mark[upd])mark[upd]=1,upd=pv[upd];
                                                                 G[eg.to][eg.back].cap+=mw;
      return upd;
     }
                                                                return {mw, dis[edd]};
    idx++;
                                                             public:
   }
                                                               void init(int n){
  }
                                                                G.clear();G.resize(n);
                                                                fa.resize(n);wh.resize(n);
 return -1;
                                                                inq.resize(n); dis.resize(n);
int Solve(int n) {
                                                               void add_edge(int st, int ed, Cap c, Wei w){
 int rt = -1, ans = 0;
                                                                G[st].emplace_back(ed,SZ(G[ed]),c,w);
 while ((rt = NegativeCycle(n)) >= 0) {
                                                                G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
 memset(mark, false, sizeof(mark));
```

```
PCW solve(int a, int b){
                                                               static int t = 0;
 ori = a, edd = b;
                                                               for (++t; u || v; swap(u, v)) {
 Cap cc=0; Wei ww=0;
                                                                if (u == 0) continue;
 while(true){
                                                                if (vis[u] == t) return u;
  PCW ret=SPFA();
                                                                vis[u] = t;
  if(ret.first==-1) break;
                                                                u = st[match[u]];
   cc+=ret.first;
                                                                if (u) u = st[pa[u]];
   ww+=ret.first * ret.second;
                                                               return 0;
 return {cc,ww};
}
                                                              void add_blossom(int u, int lca, int v) {
} mcmf;
                                                               int b = n + 1;
                                                               while (b \le n_x \& st[b]) ++b;
4.11
     Maximum Weight Graph Matching
                                                               if (b > n_x) ++n_x;
                                                               lab[b] = 0, S[b] = 0
struct WeightGraph {
static const int inf = INT_MAX;
                                                               match[b] = match[lca];
                                                               flo[b].clear();
 static const int maxn = 514;
struct edge {
                                                               flo[b].push_back(lca);
                                                               for (int x = u, y; x != lca; x = st[pa[y]])
 int u, v, w;
 edge(){}
                                                                flo[b].push_back(x), flo[b].push_back(y = st[match[x
                                                                 ]]), q_push(y);
  edge(int u, int v, int w): u(u), v(v), w(w) {}
                                                               reverse(flo[b].begin() + 1, flo[b].end())
                                                               for (int x = v, y; x != lca; x = st[pa[y]])
int n, n_x;
edge g[maxn * 2][maxn * 2];
                                                                flo[b].push_back(x), flo[b].push_back(y = st[match[x
 int lab[maxn * 2];
                                                                 ]]), q_push(y);
int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa
                                                               set_st(b, b);
                                                               for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].w
 int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
                                                                 = 0:
                                                               for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    maxn * 2];
vector<int> flo[maxn * 2];
queue<int> q;
                                                                int xs = flo[b][i];
                                                                for (int x = 1; x <= n_x; ++x)
int e_delta(const edge &e) { return lab[e.u] + lab[e.v
    ] - g[e.u][e.v].w * 2; }
                                                                 if (g[b][x].w == 0 \mid \mid e_{delta}(g[xs][x]) < e_{delta}(g[xs][x])
                                                                  [b][x]))
 void update_slack(int u, int x) { if (!slack[x] ||
    e_delta(g[u][x]) < e_delta(g[slack[x]][x])) slack[x
                                                                  g[b][x] = g[xs][x], g[x][b] = g[x][xs];
                                                                for (int x = 1; x <= n; ++x)
    ] = u; }
                                                                 if (flo_from[xs][x]) flo_from[b][x] = xs;
void set_slack(int x) {
  slack[x] = 0;
  for (int u = 1; u <= n; ++u)
                                                               set_slack(b);
  if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
   update_slack(u, x);
                                                              void expand_blossom(int b) {
                                                               for (size_t i = 0; i < flo[b].size(); ++i)
  set_st(flo[b][i], flo[b][i]);</pre>
void q_push(int x) {
 if (x \le n) q.push(x);
                                                               int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
                                                                 xr);
 else for (size_t i = 0; i < flo[x].size(); i++)</pre>
                                                               for (int i = 0; i < pr; i += 2)</pre>
    q_push(flo[x][i]);
                                                                int xs = flo[b][i], xns = flo[b][i + 1];
                                                                pa[xs] = g[xns][xs].u;
void set_st(int x, int b) {
                                                                S[xs] = 1, S[xns] = 0;
 st[x] = b;
                                                                slack[xs] = 0, set_slack(xns);
 if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)</pre>
     set_st(flo[x][i], b);
                                                                q_push(xns);
                                                               S[xr] = 1, pa[xr] = pa[b];
int get_pr(int b, int xr) {
                                                               for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
 int pr = find(flo[b].begin(), flo[b].end(), xr) - flo
 [b].begin();
if (pr % 2 == 1) {
                                                                int xs = flo[b][i];
                                                                S[xs] = -1, set_slack(xs);
   reverse(flo[b].begin() + 1, flo[b].end());
  return (int)flo[b].size() - pr;
                                                               st[b] = 0;
 }
                                                              bool on_found_edge(const edge &e) {
  return pr;
                                                               int u = st[e.u], v = st[e.v];
                                                               if (S[v] == -1) {
 void set_match(int u, int v) {
                                                                pa[v] = e.u, S[v] = 1;
 match[u] = g[u][v].v;
                                                                int nu = st[match[v]];
 if (u <= n) return;</pre>
                                                                slack[v] = slack[nu] = 0;
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u, xr);
                                                                S[nu] = 0, q_push(nu);
                                                               } else if (S[v] == 0) {
  for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo</pre>
                                                                int lca = get_lca(u, v)
    [u][i ^ 1]);
                                                                if (!lca) return augment(u,v), augment(v,u), true;
  set_match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
                                                                else add_blossom(u, lca, v);
    end());
                                                               return false;
void augment(int u, int v) {
                                                              bool matching() {
 for (; ; ) {
                                                               memset(S + 1, -1, sizeof(int) * n_x);
  int xnv = st[match[u]];
                                                               memset(slack + 1, 0, sizeof(int) * n_x);
  set_match(u, v);
                                                               q = queue<int>();
  if (!xnv) return;
                                                               for (int x = 1; x <= n_x; ++x)
   set_match(xnv, st[pa[xnv]]);
   u = st[pa[xnv]], v = xnv;
                                                                if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0,
                                                                 q_push(x);
 }
                                                               if (q.empty()) return false;
int get_lca(int u, int v) {
                                                               for (; ; ) {
```

```
while (q.size()) {
    int u = q.front(); q.pop();
    if (S[st[u]] == 1) continue;
    for (int v = 1; v <= n; ++v)
if (g[u][v].w > 0 && st[u] != st[v]) {
       if (e_delta(g[u][v]) == 0) {
       if (on_found_edge(g[u][v])) return true;
       } else update_slack(u, st[v]);
   int d = inf;
   for (int b = n + 1; b \le n_x; ++b)
    if (st[b] == b && S[b] == 1) d = min(d, lab[b] / 2)
   for (int x = 1; x <= n_x; ++x)
    if (st[x] == x && slack[x]) {
     if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x])
     else if (S[x] == 0) d = min(d, e_delta(g[slack[x
    ]][x]) / 2);
   for (int u = 1; u <= n; ++u) {
    if (S[st[u]] == 0) {
     if (lab[u] <= d) return 0;</pre>
     lab[u] -= d;
    } else if (S[st[u]] == 1) lab[u] += d;
   for (int b = n + 1; b \le n_x; ++b)
    if (st[b] == b) {
     if (S[st[b]] == 0) lab[b] += d * 2;
     else if (S[st[b]] == 1) lab[b] -= d * 2;
    }
   q = queue<int>();
   for (int x = 1; x <= n_x; ++x)
    if (st[x] == x && slack[x] && st[slack[x]] != x &&
    e_delta(g[slack[x]][x]) == 0)
     if (on_found_edge(g[slack[x]][x])) return true;
   for (int b = n + 1; b \le n_x; ++b)
    if (st[b] == b && S[b] == 1 && lab[b] == 0)
    expand_blossom(b);
  return false;
 pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n_x = n;
  int n_matches = 0;
  long long tot_weight = 0;
  for (int u = 0; u \le n; ++u) st[u] = u, flo[u].clear
    ();
  int w_max = 0;
  for (int u = 1; u <= n; ++u)
  for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);</pre>
    w_max = max(w_max, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
  while (matching()) ++n_matches;
  for (int u = 1; u <= n; ++u)</pre>
   if (match[u] && match[u] < u)</pre>
    tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
 void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
    g[vi][ui].w = wi; }
 void init(int _n) {
  n = _n;
for (int u = 1; u <= n; ++u)</pre>
   for (int v = 1; v <= n; ++v)
    g[u][v] = edge(u, v, 0);
};
5
     Math
     \lfloor \frac{n}{i} \rfloor Enumeration
T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor
5.2 Strling Number
```

5.2.1 First Kind

```
S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)
      x(x+1)\dots(x+n-1) = \sum_{k=0}^{n} S_1(n,k)x^k
     g(x) = x(x+1)\dots(x+n-1) = \sum_{k=0}^{n} a_k x^k
           \Rightarrow g(x+n) = \sum_{k=0}^{n} \frac{b_k}{(n-k)!} x^{n-k},
          b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})
```

#### 5.2.2 Second Kind

 $S_2(n,k)$  counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

# 5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 \pmod{n}
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
 if (y == 0) g=x,a=1,b=0;
 else exgcd(y,x%y,g,b,a),b=(x/y)*a;
```

# 5.4 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
 vector<T> d(output.size() + 1), me, he;
 for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
  for (size_t j = 0; j < me.size(); ++j)
d[i] += output[i - j - 2] * me[j];</pre>
  if ((d[i] -= output[i - 1]) == 0) continue;
  if (me.empty()) {
   me.resize(f = i);
   continue;
  vector<T> o(i - f - 1);
T k = -d[i] / d[f]; o.push_back(-k);
  for (T x : he) o.push_back(x * k);
  if (o.size() < me.size()) o.resize(me.size());</pre>
  for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
  if (i-f+he.size() >= me.size()) he = me, f = i;
  me = o;
 return me;
```

# Charateristic Polynomial

```
vector<vector<int>> Hessenberg(const vector<vector<int
    >> &A) {
 int N = A.size();
 vector<vector<int>> H = A;
 for (int i = 0; i < N - 2; ++i) {
  if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {
    if (H[j][i]) {
      for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
     for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k]
     ][j]);
     break:
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
 * H[i + 1][k] * (kP - coef)) % kP;
   for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
     1] + 1LL * H[k][j] * coef) % kP;
```

```
return H;
vector<int> CharacteristicPoly(const vector<vector<int
    >> &A) {
int N = A.size();
auto H = Hessenberg(A);
for (int i = 0; i < N; ++i) {
 for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
vector<vector<int>>> P(N + 1, vector<int>(N + 1));
P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
 P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
  int val = 1;
 for (int j = i - 1; j >= 0; --j) {
  int coef = 1LL * val * H[j][i - 1] % kP;
  for (int k = 0; k \le j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 }
if (N & 1) {
 for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
return P[N];
```

# 5.6 Chinese Remainder

```
1ld crt(lld ans[], lld pri[], int n){
 11d M = 1, ret = 0;
 for(int i=0;i<n;i++) M *= pri[i];</pre>
 for(int i=0;i<n;i++){</pre>
 lld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
  ret += (ans[i]*(M/pri[i])%M * iv)%M;
  ret %= M;
 return ret;
}
/*
Another:
x = a1 \% m1
x = a2 \% m2
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
0 <= x < 1cm(m1, m2)
```

### 5.7 De-Bruijn

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
  if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
} else {
 aux[t] = aux[t - p];
db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {
  aux[t] = i;
   db(t + 1, t, n, k);
}
int de_bruijn(int k, int n) {
// return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
if (k == 1) {
 res[0] = 0;
  return 1;
for (int i = 0; i < k * n; i++) aux[i] = 0;
db(1, 1, n, k);
 return sz;
```

# 5.8 DiscreteLog

```
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
  // x^? \equiv y (mod M)
  Int t = 1, c = 0, g = 1;
for (Int M_ = M; M_ > 0; M_ >>= 1)
    g = g * x % M;
  for (g = gcd(g, M); t % g != 0; ++c) {
    if (t == y) return c;
    t = t * x % M;
  if (y % g != 0) return -1;
  t /= g, y /= g, M /= g;
  Int h = 0, gs = 1;
  for (; h * h < M; ++h) gs = gs * x % M;
  unordered_map<Int, Int> bs;
  for (Int s = 0; s < h; bs[y] = ++s)
    y = y * x % M;
  for (Int s = 0; s < M; s += h) {
    t = t * gs % M;
    if (bs.count(t)) return c + s + h - bs[t];
  return -1;
}
```

#### 5.9 Extended Euler

```
a^b \equiv \begin{cases} a^b \mod \varphi(m) + \varphi(m) & \text{if } (a,m) \neq 1 \wedge b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m
```

#### 5.10 ExtendedFloorSum

```
g(a,b,c,n) = \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor
                            \left( \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \right)
                             +g(a \bmod c, b \bmod c, c, n),
                                                                                                             a \geq c \vee b \geq c
                                                                                                            n<0\vee a=0
                              \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
                            -h(c, c-b-1, a, m-1)),
h(a,b,c,n) = \sum_{i=0}^{n} \left\lfloor \frac{ai+b}{c} \right\rfloor^{2}
                            \left( \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right)
                             +\lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1)
                             +h(a \bmod c, b \bmod c, c, n)
                             +2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)
                             +2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
                                                                                                             a \geq c \vee b \geq c
                             0,
                                                                                                            n < 0 \lor a = 0
                             nm(m+1) - 2g(c, c-b-1, a, m-1)
                            -2f(c, c-b-1, a, m-1) - f(a, b, c, n), otherwise
```

#### 5.11 Fast Fourier Transform

```
const int mod = 1000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
  static_assert (M1 <= M2 && M2 <= M3);</pre>
  constexpr int64_t r12 = modpow(M1, M2-2, M2);
  constexpr int64_t r13 = modpow(M1, M3-2, M3);
constexpr int64_t r23 = modpow(M2, M3-2, M3);
  constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
  B = (B - A + M2) * r12 % M2;
  C = (C - A + M3) * r13 % M3;
  C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
 for (int i = 0; i <= maxn; i++)</pre>
  omega[i] = cplx(cos(2 * pi * j / maxn),
     sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
 int z = __builtin_ctz(n) - 1;
 for (int i = 0; i < n; ++i) {
```

```
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
  int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^=(i >> j & 1)<<(z - j);
  if (x > i) swap(v[x], v[i]);
                                                               // @return sum_{i=0^{n-1} floor((ai + b)/m) mod 2^64
                                                               1lu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
for (int s = 2; s <= n; s <<= 1) {
                                                                llu ans = 0:
  int z = s >> 1;
                                                                 while (true)
  for (int i = 0; i < n; i += s) {
                                                                  if (a >= m) {
                                                                   ans += n * (n - 1) / 2 * (a / m); a %= m;
  for (int k = 0; k < z; ++k) {
   cplx x = v[i + z + k] * omega[maxn / s * k];
   v[i+z+k] = v[i+k] - x;
                                                                  if (b >= m) {
                                                                  ans += n * (b / m); b %= m:
    v[i+k] = v[i+k] + x;
  }
                                                                  llu y_max = a * n + b;
                                                                  if (y_max < m) break;</pre>
                                                                  // y_max < m * (n + 1)
void ifft(vector<cplx> &v, int n) {
                                                                  // floor(y_max / m) <= n
fft(v, n); reverse(v.begin() + 1, v.end());
                                                                  n = (11u)(y_max / m), b = (11u)(y_max % m);
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);
                                                                  swap(m, a);
VL convolution(const VI &a, const VI &b) {
                                                                 return ans;
// Should be able to handle N <= 10^5, C <= 10^4
                                                               11d floor_sum(lld n, lld m, lld a, lld b) {
int sz = 1;
while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
                                                                 llu ans = 0;
vector<cplx> v(sz);
                                                                 if (a < 0) {
for (int i = 0; i < sz; ++i)
                                                                 11u \ a2 = (a \% m + m) \% m;
 double re = i < a.size() ? a[i] : 0;</pre>
                                                                  ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
 double im = i < b.size() ? b[i] : 0;</pre>
                                                                  a = a2:
  v[i] = cplx(re, im);
                                                                 if (b < 0) {
fft(v, sz);
                                                                 11u b2 = (b \% m + m) \% m;
for (int i = 0; i <= sz / 2; ++i) {
                                                                  ans -= 1ULL * n * ((b2 - b) / m);
 int j = (sz - i) & (sz - 1);
                                                                 b = b2:
 cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
    * cplx(0, -0.25);
                                                                 return ans + floor_sum_unsigned(n, m, a, b);
 if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
    ].conj()) * cplx(0, -0.25);
                                                               5.13
                                                                      FWT
  v[i] = x;
                                                               /* xor convolution:
                                                                * x = (x0, x1) , y = (y0, y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
 ifft(v, sz);
VL c(sz);
for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
                                                                * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1)
* z = (1/2) * z''
VI convolution_mod(const VI &a, const VI &b, int p) {
int sz = 1;
                                                                 * or convolution:
while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
                                                                 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
vector<cplx> fa(sz), fb(sz);
                                                                 * and convolution:
                                                                 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
for (int i = 0; i < (int)a.size(); ++i)</pre>
 fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
                                                               const LL MOD = 1e9+7;
for (int i = 0; i < (int)b.size(); ++i)</pre>
                                                               inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
 fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
                                                                 for( int d = 1 ; d < N ; d <<= 1 ) {</pre>
 fft(fa, sz), fft(fb, sz);
                                                                  int d2 = d << 1;
double r = 0.25 / sz;
                                                                  for( int s = 0 ; s < N ; s += d2 )
                                                                  for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
    LL ta = x[ i ] , tb = x[ j ];
    x[ i ] = ta+tb;
cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1); for (int i = 0; i <= (sz >> 1); ++i) {
 int j = (sz - i) & (sz - 1);
  cplx a1 = (fa[i] + fa[j].conj());
                                                                    x[ j ] = ta-tb;
                                                                    if( x[ i ] >= MOD ) x[ i ] -= MOD;
if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
 cplx a2 = (fa[i] - fa[j].conj()) * r2;
 cplx b1 = (fb[i] + fb[j].conj()) * r3;
  cplx b2 = (fb[i] - fb[j].conj()) * r4;
  if (i != j) {
   cplx c1 = (fa[j] + fa[i].conj());
                                                                 if( inv )
   cplx c2 = (fa[j] - fa[i].conj()) * r2;
                                                                 for( int i = 0 ; i < N ; i++ ) {</pre>
                                                                  x[ i ] *= inv( N, MOD );
x[ i ] %= MOD;
  cplx d1 = (fb[j] + fb[i].conj()) * r3;
   cplx d2 = (fb[j] - fb[i].conj()) * r4;
   fa[i] = c1 * d1 + c2 * d2 * r5;
   fb[i] = c1 * d2 + c2 * d1;
                                                               }
                                                                      Miller Rabin
  fa[j] = a1 * b1 + a2 * b2 * r5;
 fb[j] = a1 * b2 + a2 * b1;
                                                               bool isprime(llu x) {
                                                                 static auto witn = [](llu a, llu u, llu n, int t) {
fft(fa, sz), fft(fb, sz);
                                                                 if (!a) return false;
vector<int> res(sz);
                                                                  while (t--) {
                                                                   llu a2 = mmul(a, a, n);
 for (int i = 0; i < sz; ++i) {
 long long a = round(fa[i].re), b = round(fb[i].re),
                                                                   if (a2 == 1 && a != 1 && a != n - 1) return true;
       c = round(fa[i].im);
 res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
                                                                 }
                                                                  return a != 1;
return res;
}}
                                                                if (x < 2) return false;</pre>
                                                                 if (!(x & 1)) return x == 2;
```

int t = \_\_builtin\_ctzll(x - 1);

```
11u \text{ odd} = (x - 1) >> t;
                                                                  if(p * i >= N) break;
                                                                 sieved[p * i] = true;
 for (llu m:
  {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
                                                                 if(p % i == 0) break;
  if (witn(mpow(m % x, odd, x), odd, x, t))
  return false:
 return true;
                                                              11d phi(11d m, 11d n) {
                                                               static constexpr int MM = 80000, NN = 500;
5.15 NTT
                                                               static lld val[MM][NN];
                                                               if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
template <int mod, int G, int maxn>
                                                               if(n == 0) return m;
struct NTT {
                                                               if(primes[n] >= m) return 1;
 static_assert (maxn == (maxn & -maxn));
                                                               11d ret = phi(m,n-1)-phi(m/primes[n],n-1);
 int roots[maxn];
                                                               if(m < MM\&n < NN) val[m][n] = ret+1;
 NTT () {
                                                               return ret:
  int r = modpow(G, (mod - 1) / maxn);
  for (int i = maxn >> 1; i; i >>= 1) {
                                                              1ld pi_count(1ld);
   roots[i] = 1;
                                                              11d P2(11d m, 11d n) {
   for (int j = 1; j < i; j++)
roots[i + j] = modmul(roots[i + j - 1], r);</pre>
                                                               11d sm = square_root(m), ret = 0;
                                                               for(lld i = n+1;primes[i]<=sm;i++)</pre>
   r = modmul(r, r);
                                                                ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
                                                               return ret;
 // n must be 2^k, and 0 \le F[i] < mod
                                                              11d pi_count(11d m) {
 void operator()(int F[], int n, bool inv = false) {
                                                               if(m < N) return pi[m];</pre>
  for (int i = 0, j = 0; i < n; i++) {
  if (i < j) swap(F[i], F[j]);</pre>
                                                               11d n = pi_count(cube_root(m));
                                                               return phi(m, n) + n - 1 - P2(m, n);
   for (int k = n > 1; (j^k < k; k > = 1);
  for (int s = 1; s < n; s *= 2) {
                                                              5.18 Pollard Rho
   for (int i = 0; i < n; i += s * 2) {
                                                              // does not work when n is prime
    for (int j = 0; j < s; j++) {
                                                              // return any non-trivial factor
     int a = F[i+j];
                                                              llu pollard_rho(llu n) {
     int b = modmul(F[i+j+s], roots[s+j]);
                                                               static auto f = [](llu x, llu k, llu m) {
     F[i+j] = modadd(a, b); // a + b
                                                                   return add(k, mul(x, x, m), m); };
     F[i+j+s] = modsub(a, b); // a - b
                                                               if (!(n & 1)) return 2;
                                                               mt19937 rnd(120821011);
   }
                                                               while (true) {
                                                                llu y = 2, yy = y, x = rnd() % n, t = 1;
for (llu sz = 2; t == 1; sz <<= 1, y = yy) {
  if (inv) {
   int invn = modinv(n);
                                                                 for (llu i = 0; t == 1 && i < sz; ++i) {
   for (int i = 0; i < n; i++)</pre>
                                                                  yy = f(yy, x, n);
   F[i] = modmul(F[i], invn);
                                                                  t = gcd(yy > y ? yy - y : y - yy, n);
   reverse(F + 1, F + n);
 }
                                                                if (t != 1 && t != n) return t;
NTT<2013265921, 31, 1048576> ntt;
5.16 Partition Number
                                                              5.19 Polynomial Operations
int b = sqrt(n);
                                                              using V = vector<int>
ans[0] = tmp[0] = 1;
                                                              #define fi(1, r) for (int i = int(1); i < int(r); ++i)
for (int i = 1; i <= b; i++) {
                                                              template <int mod, int G, int maxn> struct Poly : V {
 for (int rep = 0; rep < 2; rep++)</pre>
                                                               static uint32_t n2k(uint32_t n) {
 for (int j = i; j <= n - i * i; j++)
                                                                if (n <= 1) return 1;</pre>
 modadd(tmp[j], tmp[j-i]);
for (int j = i * i; j <= n; j++)</pre>
                                                                return 1u << (32 - __builtin_clz(n - 1));</pre>
 modadd(ans[j], tmp[j - i * i]);
                                                               static NTT<mod,G,maxn> ntt; // coefficients in [0, P)
                                                               explicit Poly(int n = 1) : V(n) {}
                                                               Poly(const V &v) : V(v) {}
5.17 Pi Count (Linear Sieve)
                                                               Poly(const Poly &p, size_t n) : V(n) {
static constexpr int N = 1000000 + 5;
                                                                copy_n(p.data(), min(p.size(), n), data());
11d pi[N];
                                                               Poly &irev() { return reverse(data(), data() + size())
vector<int> primes;
bool sieved[N];
                                                                     *this; }
11d cube_root(11d x){
                                                               Poly &isz(int sz) { return resize(sz), *this; }
                                                               Poly aidd(const\ Poly\ arhs) { // n() == rhs.n()}
 1ld s=cbrt(x-static_cast<long double>(0.1));
                                                                fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
 while(s*s*s <= x) ++s;
                                                                return *this:
return s-1;
11d square_root(11d x){
                                                               Poly &imul(int k) {
1ld s=sqrt(x-static_cast<long double>(0.1));
                                                                fi(0, size())(*this)[i] = modmul((*this)[i], k);
 while(s*s <= x) ++s;</pre>
                                                                return *this;
 return s-1;
                                                               Poly Mul(const Poly &rhs) const {
void init(){
                                                                const int sz = n2k(size() + rhs.size() - 1);
                                                                Poly X(*this, sz), Y(rhs, sz);
ntt(X.data(), sz), ntt(Y.data(),
primes.reserve(N)
 primes.push_back(1);
 for(int i=2;i<N;i++) {</pre>
                                                                fi(0, sz) X[i] = modmul(X[i], Y[i]);
  if(!sieved[i]) primes.push_back(i);
                                                                ntt(X.data(), sz, true);
  pi[i] = !sieved[i] + pi[i-1];
                                                                return X.isz(size() + rhs.size() - 1);
  for(int p: primes) if(p > 1) {
```

```
Poly Inv() const { // coef[0] != 0
 if (size() == 1) return V{modinv(*begin())};
                                                                static int LinearRecursion(const V &a, const V &c,
 const int sz = n2k(size() * 2);
                                                                   int64_t n) { // a_n = \sum c_j a_(n-j)}
 Poly X = Poly(*this, (size() + 1) / 2).Inv().isz(sz),
                                                                 const int k = (int)a.size();
    Y(*this, sz);
                                                                 assert((int)c.size() == k + 1);
                                                                 Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
 ntt(X.data(), sz), ntt(Y.data(), sz);
 fi(0, sz) X[i] = modmul(X[i], modsub(2, modmul(X[i],
                                                                 fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
   Y[i])));
                                                                 Clkl = 1:
                                                                 while (n) {
 ntt(X.data(), sz, true);
                                                                  if (n % 2) W = W.Mul(M).DivMod(C).second;
 return X.isz(size());
                                                                  n /= 2, M = M.Mul(M).DivMod(C).second;
Poly Sqrt() const { // coef[0] \in [1, mod)^2
 if (size() == 1) return V{QuadraticResidue((*this)
                                                                 int ret = 0:
                                                                 fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
   [0], mod)};
 Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
                                                                 return ret;
   size());
 return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
                                                               #undef fi
   + 1):
                                                               using Poly_t = Poly<998244353, 3, 1 << 20>;
pair<Poly, Poly> DivMod(const Poly &rhs) const {
                                                               template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
if (size() < rhs.size()) return {V{0}, *this};</pre>
                                                               5.20 Quadratic residue
 const int sz = size() - rhs.size() + 1;
 Poly X(rhs); X.irev().isz(sz);
                                                               struct S {
 Poly Y(*this); Y.irev().isz(sz);
                                                                int MOD, w;
 Poly Q = Y.Mul(X.Inv()).isz(sz).irev();
                                                                int64_t x, y;
 X = rhs.Mul(Q), Y = *this;
fi(0, size()) Y[i] = modsub(Y[i], X[i]);
                                                                S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
                                                                 : MOD(m), w(w_{-}), x(x_{-}), y(y_{-}) {}
 return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
                                                                S operator*(const S &rhs) const {
                                                                 int w_{-} = w;
Poly Dx() const {
                                                                 if (w<sub>_</sub> == -1) w<sub>_</sub> = rhs.w;
                                                                 assert(w_! = -1 \text{ and } w_ == rhs.w);
 Poly ret(size() - 1);
                                                                 return { MOD, w_,
(x * rhs.x + y * rhs.y % MOD * w) % MOD,
(x * rhs.y + y * rhs.x) % MOD };
 fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
   1]);
 return ret.isz(max<int>(1, ret.size()));
Poly Sx() const {
                                                               };
Poly ret(size() + 1);
                                                               int get_root(int n, int P) {
 fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
                                                                 if (P == 2 or n == 0) return n;
                                                                 if (qpow(n, (P - 1) / 2, P) != 1) return -1;
auto check = [&](int x) {
   this)[i]);
 return ret:
                                                                 return qpow(x, (P - 1) / 2, P); };
if (check(n) == P-1) return -1;
Poly Ln() const { // coef[0] == 1
  return Dx().Mul(Inv()).Sx().isz(size());
                                                                 int64_t a; int w; mt19937 rnd(7122);
                                                                 do { a = rnd() % P;
                                                                   w = ((a * a - n) % P + P) % P;
Poly Exp() const \{ // coef[0] == 0 \}
 if (size() == 1) return V{1};
                                                                 } while (check(w) != P - 1);
 Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size
                                                                 return qpow(S(P, w, a, 1), (P + 1) / 2).x;
   ());
 Poly Y = X.Ln(); Y[0] = mod - 1;
                                                               5.21 Simplex
 fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
 return X.Mul(Y).isz(size());
                                                               namespace simplex {
                                                               // maximize c^Tx under Ax <= B
                                                               // return VD(n, -inf) if the solution doesn't exist
Poly Pow(const string &K) const {
 int nz = 0;
                                                               // return VD(n, +inf) if the solution is unbounded
 while (nz < size() && !(*this)[nz]) ++nz;</pre>
                                                               using VD = vector<double>;
 int nk = 0, nk2 = 0;
                                                               using VVD = vector<vector<double>>;
 for (char c : K) {
                                                               const double eps = 1e-9;
  nk = (nk * 10 + c - '0') % mod;
                                                               const double inf = 1e+9;
  nk2 = nk2 * 10 + c - '0'
                                                               int n, m;
  if (nk2 * nz >= size())
                                                               VVD d;
   return Poly(size());
                                                               vector<int> p, q;
  nk2 %= mod - 1;
                                                               void pivot(int r, int s) {
                                                                double inv = 1.0 / d[r][s];
 if (!nk && !nk2) return Poly(V{1}, size());
                                                                for (int i = 0; i < m + 2; ++i)
 Poly X = V(data() + nz, data() + size() - nz * (nk2 -
                                                                 for (int j = 0; j < n + 2; ++j)
    1));
                                                                  if (i != r && j != s)
 int x0 = X[0];
                                                                   d[i][j] -= d[r][j] * d[i][s] * inv;
                                                                for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
 return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
   modpow(x0, nk2)).irev().isz(size()).irev();
                                                                d[r][s] = inv; swap(p[r], q[s]);
Poly InvMod(int L) { // (to evaluate linear recursion)
Poly R\{1, 0\}; // *this * R mod x^L = 1 (*this[0] ==
                                                               bool phase(int z) {
                                                                int x = m + z;
 for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                                while (true) {
  Poly 0 = R.Mul(Poly(data(), min<int>(2 << level,
                                                                 int s = -1;
   size())));
                                                                 for (int i = 0; i <= n; ++i) {</pre>
                                                                  if (!z && q[i] == -1) continue;
if (s == -1 || d[x][i] < d[x][s]) s = i;</pre>
  Poly Q(2 << level); Q[0] = 1;
for (int j = (1 << level); j < (2 << level); ++j)
   Q[j] = modsub(mod, O[j]);
  R = R.Mul(Q).isz(4 << level);
                                                                 if (d[x][s] > -eps) return true;
                                                                 int r = -1;
                                                                 for (int i = 0; i < m; ++i) {
 return R.isz(L);
```

```
if (d[i][s] < eps) continue;</pre>
    if (r == -1 ||
     d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
  pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
 m = b.size(), n = c.size();
d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)
  p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
 for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];
 q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
 for (int i = 1; i < m;</pre>
  if (d[i][n + 1] < d[r][n + 1]) r = i;
 if (d[r][n + 1] < -eps) {</pre>
  pivot(r, n);
  if (!phase(1) || d[m + 1][n + 1] < -eps)</pre>
    return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
   int s = min_element(d[i].begin(), d[i].end() - 1)
         - d[i].begin();
   pivot(i, s);
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)
  if (p[i] < n) x[p[i]] = d[i][n + 1];
 return x;
}}
5.22
       Simplex Construction
Standard form: maximize \sum_{1 \le i \le n} c_i x_i such that for all 1 \le j \le m,
\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j. \text{and } x_i \geq 0 \text{ for all } 1 \leq i \leq n.
  1. In case of minimization, let c'_i = -c_i
  2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
  3. \sum_{1 < i < n} A_{ji} x_i = b_j
         • \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
         • \sum_{1 \le i \le n} A_{ji} x_i \ge b_j
  4. If x_i has no lower bound, replace x_i with x_i - x_i'
```

# 6 Geometry

# 6.1 Basic Geometry

```
#define IM imag
#define RE real
using lld = int64_t;
using llf = long double;
using PT = std::complex<1ld>;
using PTF = std::complex<llf>;
auto toPTF(PT p) { return PTF{RE(p), IM(p)}; }
int sgn(lld x) { return (x > 0) - (x < 0); }</pre>
11d dot(PT a, PT b) { return RE(conj(a) * b); }
11d cross(PT a, PT b) { return IM(conj(a) * b); }
int ori(PT a, PT b, PT c) {
return sgn(cross(b - a, c - a));
bool operator<(const PT &a, const PT &b)
return RE(a) != RE(b) ? RE(a) < RE(b) : IM(a) < IM(b);
int quad(PT p) {
 return (IM(p) == 0) // use sgn for PTF
  ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
int argCmp(PT a, PT b) {
// -1 / 0 / 1 <-> < / == / > (atan2)
 int qa = quad(a), qb = quad(b);
 if (qa != qb) return sgn(qa - qb);
 return sgn(cross(b, a));
```

```
template <typename V> llf area(const V & pt) {
 11d ret = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 return ret / 2.0;
PT rot90(PT p) { return PT{-IM(p), RE(p)}; }
PTF project(PTF p, PTF q) { // p onto q
return dot(p, q) * q / dot(q, q);
11f FMOD(11f x) {
 if (x < -PI) x += PI * 2;
 if (x > PI) x -= PI * 2;
 return x;
6.2 Segment & Line Intersection
struct Segment {
 PT st, dir; // represent st + t*dir for 0<=t<=1
 Segment(PT s, PT e) : st(s), dir(e - s) {}
 static bool valid(lld p, lld q) {
  // is there t s.t. 0 <= t <= 1 && qt == p ?
  if (q < 0) q = -q, p = -p;
  return 0 <= p && p <= q;
bool isInter(Segment A, PT P) {
 if (A.dir == PT(0)) return P == A.st;
 return cross(P - A.st, A.dir) == 0 &&
  Segment::valid(dot(P - A.st, A.dir), norm(A.dir));
template <typename U, typename V>
bool isInter(U A, V B) {
  if (cross(A.dir, B.dir) == 0)
  return // handle parallel yourself
   isInter(A, B.st) || isInter(A, B.st+B.dir) ||
   isInter(B, A.st) || isInter(B, A.st+A.dir);
 PT D = B.st - A.st;
 11d C = cross(A.dir, B.dir)
 return U::valid(cross(D, A.dir), C) &&
   V::valid(cross(D, B.dir), C);
struct Line {
 PT st, ed, dir;
 Line (PT s, PT e)
  : st(s), ed(e), dir(e - s) {}
PTF intersect(const Line &A, const Line &B) {
 11f t = cross(B.st - A.st, B.dir) /
  llf(cross(A.dir, B.dir))
 return toPTF(A.st) + PTF(t) * toPTF(A.dir);
6.3 2D Convex Hull
// returns a convex hull in counterclockwise order
// for a non-strict one, change cross >= to >
vector<PT> convex_hull(vector<PT> p) {
 sort(all(p));
 if (p[0] == p.back()) return {p[0]};
 int n = p.size(), t = 0;
 vector<PT> h(n + 1);
 for (int _{-} = 2, s = 0; _{--}; s = --t, reverse(all(p)))
  for (PT i : p) {
   while (t > s + 1 \&\& cross(i, h[t-1], h[t-2]) >= 0)
    t --
   h[t++] = i;
 return h.resize(t), h;
6.4 3D Convex Hull
// return the faces with pt indexes
int flag[MXN][MXN];
struct Point{
 ld x,y,z;
 Point operator * (const ld &b) const {
  return (Point) {x*b, y*b, z*b};}
 Point operator * (const Point &b) const {
```

return(Point){y\*b.z-b.y\*z,z\*b.x-b.z\*x,x\*b.y-b.x\*y};

```
Point ver(Point a, Point b, Point c) {
  return (b - a) * (c - a);}
                                                                     }
                                                                   }
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0;
                                                                  if (found) rebuild_m(i + 1);
 REP(i,n) REP(j,n) flag[i][j] = 0;
                                                                  else m[kx][ky][kz] = i;
 vector<Face> now;
 now.emplace_back(0,1,2);
                                                                 return d:
 now.emplace_back(2,1,0);
 for (int i=3; i<n; i++){</pre>
                                                                     Simulated Annealing
                                                                6.7
  ftop++; vector<Face> next;
  REP(j, SZ(now)) {
                                                                11f anneal() {
   Face& f=now[j]; int ff = 0;
ld d=(pt[i]-pt[f.a]).dot(
                                                                 mt19937 rnd_engine( seed );
                                                                 uniform_real_distribution< llf > rnd( 0, 1 );
     ver(pt[f.a], pt[f.b], pt[f.c]));
                                                                 const 11f dT = 0.001;
   if (d <= 0) next.push_back(f);</pre>
                                                                  // Argument p
   if (d > 0) ff=ftop;
                                                                 11f S_cur = calc( p ), S_best = S_cur;
                                                                 for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
   else if (d < 0) ff=-ftop;
   flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
                                                                  // Modify p to p_prime
const llf S_prime = calc( p_prime );
  REP(j, SZ(now)) {
                                                                  const 11f delta_c = S_prime - S_cur;
   Face& f=now[j]
                                                                  11f prob = min( ( 11f ) 1, exp( -delta_c / T ) );
                                                                  if ( rnd( rnd_engine ) <= prob )</pre>
   if (flag[f.a][f.b] > 0 &&
                                                                   S_cur = S_prime, p = p_prime;
     flag[f.a][f.b] != flag[f.b][f.a])
    next.emplace_back(f.a,f.b,i);
                                                                  if ( S_prime < S_best ) // find min</pre>
   if (flag[f.b][f.c] > 0 &&
   flag[f.b][f.c] != flag[f.c][f.b])
                                                                   S_best = S_prime, p_best = p_prime;
    next.emplace_back(f.b,f.c,i);
                                                                 return S_best;
   if (flag[f.c][f.a] > 0 &&
  flag[f.c][f.a] != flag[f.a][f.c])
                                                                6.8 Half Plane Intersection
    next.emplace_back(f.c,f.a,i);
                                                                // cross(pt-line.st, line.dir)<=0 <-> pt in half plane
  now=next;
                                                                bool operator<(const Line &lhs, const Line &rhs) {</pre>
 }
                                                                  if (int cmp = argCmp(lhs.dir, rhs.dir))
 return now;
                                                                     return cmp == -1;
                                                                   return ori(lhs.st, lhs.ed, rhs.st) < 0;
6.5 2D Farthest Pair
// stk is from convex hull
                                                                // intersect function is in "Segment Intersect"
n = (int)(stk.size());
                                                                llf HPI(vector<Line> &lines) {
int pos = 1, ans = 0; stk.push_back(stk[0]);
                                                                  sort(lines.begin(), lines.end());
for(int i=0;i<n;i++) {</pre>
                                                                  deque<Line> que;
 while(abs(cross(stk[i+1]-stk[i],
                                                                  deque<PTF> pt;
   stk[(pos+1)%n]-stk[i])) >
                                                                  que.push_back(lines[0]);
   abs(cross(stk[i+1]-stk[i],
                                                                  for (int i = 1; i < (int)lines.size(); i++) {</pre>
 stk[pos]-stk[i]))) pos = (pos+1)%n;
ans = max({ans, dis(stk[i], stk[pos])},
                                                                     if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
                                                                      continue:
                                                                #define POP(L, R) \
  dis(stk[i+1], stk[pos])});
                                                                     while (pt.size() > 0 \
                                                                       && ori(L.st, L.ed, pt.back()) < 0) \
6.6 kD Closest Pair (3D ver.)
                                                                       pt.pop_back(), que.pop_back(); \
11f solve(vector<P> v) {
                                                                     while (pt.size() > 0 \
 shuffle(v.begin(), v.end(), mt19937());
                                                                       && ori(R.st, R.ed, pt.front()) < 0) \
 unordered_map<lld, unordered_map<lld,
                                                                       pt.pop_front(), que.pop_front();
  unordered_map<lld, int>>> m;
                                                                     POP(lines[i], lines[i])
 llf d = dis(v[0], v[1]);
                                                                     pt.push_back(intersect(que.back(), lines[i]));
 auto Idx = [&d] (11f x) -> 11d {
                                                                     que.push_back(lines[i]);
  return round(x \star 2 / d) + 0.1; };
 auto rebuild_m = [&m, &v, &Idx](int k) {
                                                                  POP(que.front(), que.back())
  m.clear();
                                                                  if (que.size() <= 1 ||</pre>
  for (int i = 0; i < k; ++i)
m[Idx(v[i].x)][Idx(v[i].y)]</pre>
                                                                     argCmp(que.front().dir, que.back().dir) == 0)
                                                                     return 0
    [Idx(v[i].z)] = i;
                                                                  pt.push_back(intersect(que.front(), que.back()));
 }; rebuild_m(2);
                                                                   return area(pt);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
                                                                6.9 Minkowski Sum
     kz = Idx(v[i].z); bool found = false;
  for (int dx = -2; dx <= 2; ++dx) {
                                                                vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
   const 11d nx = dx + kx;
                                                                 hull(A), hull(B);
                                                                 vector<pll> C(1, A[0] + B[0]), s1, s2;
for(int i = 0; i < SZ(A); ++i)
s1.pb(A[(i + 1) % SZ(A)] - A[i]);</pre>
   if (m.find(nx) == m.end()) continue;
   auto& mm = m[nx];
for (int dy = -2; dy <= 2; ++dy) {</pre>
    const 11d ny = dy + ky;
                                                                 for(int i = 0; i < SZ(B); i++)</pre>
    if (mm.find(ny) == mm.end()) continue;
                                                                  s2.pb(B[(i + 1) % SZ(B)] - B[i]);
                                                                 for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
    auto& mmm = mm[ny];
                                                                  if (p2 >= SZ(B)
    for (int dz = -2; dz <= 2; ++dz) {
     const 1ld nz = dz + kz;
                                                                     || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
     if (mmm.find(nz) == mmm.end()) continue;
                                                                   C.pb(C.back() + s1[p1++]);
     const int p = mmm[nz];
                                                                  else
                                                                   C.pb(C.back() + s2[p2++]);
     if (dis(v[p], v[i]) < d) {</pre>
      d = dis(v[p], v[i]);
                                                                 return hull(C), C;
      found = true;
```

#### 6.10 Circle Class

```
struct Circle { PTF o; llf r; };
vector<llf> intersectAngle(Circle A, Circle B) {
PTF dir = B.o - A.o; 11f d2 = norm(dir);
if (norm(A.r - B.r) >= d2) // norm(x) := |x|^2
 if (A.r < B.r) return {-PI, PI}; // A in B</pre>
 else return {}; // B in A
 if (norm(A.r + B.r) <= d2) return {};</pre>
llf dis = abs(dir), theta = arg(dir);
llf phi = acos((A.r * A.r + d2 - B.r * B.r) /
   (2 * A.r * dis));
11f L = FMOD(theta - phi), R = FMOD(theta + phi);
return { L, R };
vector<PTF> intersectPoint(Circle a, Circle b) {
11f d = abs(a.o - b.o);
if (d >= b.r+a.r || d <= abs(b.r-a.r)) return {};</pre>
11f dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
PTF dir = (a.o - b.o) / d;
PTF u = dir*d1 + b.o;
PTF v = rot90(dir) * sqrt(max<llf>(0, b.r*b.r-d1*d1));
return {u + v, u - v};
```

### 6.11 Intersection of line and Circle

```
vector<PTF> line_interCircle(const PTF &p1,
   const PTF &p2, const PTF &c, const double r) {
  PTF ft = p1 + project(c-p1, p2-p1), vec = p2-p1;
  llf dis = abs(c - ft);
  if (abs(dis - r) < eps) return {ft};
  if (dis > r) return {};
  vec = vec * sqrt(r * r - dis * dis) / abs(vec);
  return {ft + vec, ft - vec};
}
```

## 6.12 Intersection of Polygon and Circle

```
// Divides into multiple triangle, and sum up
  test by HDU2892
11f _area(PTF pa, PTF pb, llf r) {
if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
 if (abs(pb) < eps) return 0;</pre>
11f S, h, theta;
llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
11f cosB = dot(pb, pb - pa) / a / c, B = acos(cosB);
11f cosC = dot(pa, pb) / a / b, C = acos(cosC);
if (a > r) {
 S = (C / 2) * r * r;

h = a * b * sin(C) / c;
 if (h < r && B < PI / 2)
   S = (acos(h / r) * r * r - h * sqrt(r*r - h*h));
 } else if (b > r) {
 theta = \overrightarrow{PI} - \overrightarrow{B} - \overrightarrow{a}sin(sin(B) / r * a);
 S = 0.5 * a*r * sin(theta) + (C - theta) / 2 * r*r;
} else
 S = 0.5 * sin(C) * a * b;
 return S;
llf area_poly_circle(const vector<PTF> &poly,
  const PTF &0, const llf r) {
 11f S = 0;
 for (int i = 0, N = poly.size(); i < N; ++i)</pre>
 S += _area(poly[i] - 0, poly[(i + 1) % N] - 0, r) *
     ori(0, poly[i], poly[(i + 1) % N]);
return fabs(S);
```

# 6.13 Point & Hulls Tangent

```
#define above(P, Vi, Vj) (ori(P, Vi, Vj) > 0) // true
    if Vi is above Vj
#define below(P, Vi, Vj) (ori(P, Vi, Vj) < 0) // true
    if Vi is below Vj
// Rtangent_PointPolyC(): binary search for convex
    polygon right tangent
// Input: P = a 2D point (exterior to the polygon)
// n = number of polygon vertices
// V = array of vertices for a 2D convex polygon
    with V[n] = V[0]
// Return: index "i" of rightmost tangent point V[i]
int Rtangent_PointPolyC(PT P, int n, PT *V) {</pre>
```

```
int a, b, c
 int upA, dnC;
 if (below(P, V[1], V[0]) && !above(P, V[n - 1], V[0]))
  return 0:
 for (a = 0, b = n;;) {
  c = (a + b) / 2;
  dnC = below(P, V[c + 1], V[c]);
  if (dnC && !above(P, V[c - 1], V[c]))
   return c;
  upA = above(P, V[a + 1], V[a]);
  if (upA) {
   if (dnC) {
    b = c;
   } else {
    if (above(P, V[a], V[c]))
     b = c;
    else
     a = c:
   }
  } else {
   if (!dnC) {
    a = c;
   } else {
    if (below(P, V[a], V[c]))
     b = c;
    else
     a = c;
}
// Ltangent_PointPolyC(): binary search for convex
    polygon left tangent
    Input: P = a 2D point (exterior to the polygon)
//
//
        n = number of polygon vertices
        V = array of vertices for a 2D convex polygon
//
    with V[n]=V[0]
    Return: index "i" of leftmost tangent point V[i]
int Ltangent_PointPolyC(PT P, int n, PT *V) {
 int a, b, c;
 int dnA, dnC;
 if (above(P, V[n - 1], V[0]) && !below(P, V[1], V[0]))
  return 0;
 for (a = 0, b = n;;) {
  c = (a + b) / 2;
  dnC = below(P, V[c + 1], V[c]);
  if (above(P, V[c - 1], V[c]) && !dnC)
   return c
  dnA = below(P, V[a + 1], V[a]);
  if (dnA) {
   if (!dnC) {
    b = c;
   } else {
    if (below(P, V[a], V[c]))
     b = c;
    else
  } else {
   if (dnC) {
    a = c;
   } else {
    if (above(P, V[a], V[c]))
     b = c;
    else
     a = c;
}
```

# 6.14 Convex Hulls Tangent

```
// RLtangent_PolyPolyC(): get the RL tangent between
   two convex polygons
// Input: m = number of vertices in polygon 1
```

```
V = array of vertices for convex polygon 1 with
                                                                  c.r = dist(pts[i], c.o);
     V[m]=V[0]
                                                                  for (int k = 0; k < j; k++) {
        n = number of vertices in polygon 2
                                                                   if (dist(pts[k], c.o) <= c.r) continue;</pre>
//
        W = array of vertices for convex polygon 2 with
                                                                   c = getCircum(pts[i], pts[j], pts[k]);
     W[n]=W[0]
   Output: *t1 = index of tangent point V[t1] for
                                                                 }
    polygon 1
        *t2 = index of tangent point W[t2] for polygon
//
                                                                return c;
void RLtangent_PolyPolyC(int m, PT *V, int n, PT *W,
                                                                     KDTree (Nearest Point)
                                                               6.17
    int *t1, int *t2) {
int ix1, ix2; // search indices for polygons 1 and 2
                                                              const int MXN = 100005;
                                                               struct KDTree {
 // first get the initial vertex on each polygon
                                                                struct Node {
ix1 = Rtangent_PointPolyC(W[0], m, V); // right
                                                                 int x,y,x1,y1,x2,y2;
    tangent from W[0] to V
                                                                 int id,f;
Node *L, *R;
ix2 = Ltangent_PointPolyC(V[ix1], n, W); // left
    tangent from V[ix1] to W
                                                                } tree[MXN], *root;
                                                                int n;
                                                                LL dis2(int x1, int y1, int x2, int y2) {
LL dx = x1-x2, dy = y1-y2;
 // ping-pong linear search until it stabilizes
int done = false; // flag when done
while (done == false) {
                                                                 return dx*dx+dy*dy;
 done = true; // assume done until..
 while (ori(W[ix2], V[ix1], V[ix1 + 1]) <= 0) {</pre>
                                                                static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
  ++ix1; // get Rtangent from W[ix2] to V
                                                                static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
                                                                void init(vector<pair<int,int>> ip) {
 while (ori(V[ix1], W[ix2], W[ix2 - 1]) >= 0) {
                                                                 n = ip.size();
             // get Ltangent from V[ix1] to W
                                                                 for (int i=0; i<n; i++) {</pre>
   done = false; // not done if had to adjust this
                                                                  tree[i].id = i;
                                                                  tree[i].x = ip[i].first;
  }
                                                                  tree[i].y = ip[i].second;
 *t1 = ix1;
*t2 = ix2;
                                                                 root = build_tree(0, n-1, 0);
return;
                                                                Node* build_tree(int L, int R, int d) {
                                                                 if (L>R) return nullptr
6.15
      Tangent line of Two Circle
                                                                 int M = (L+R)/2; tree[M].f = d%2;
vector<Line>
                                                                 nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
tanline(const Circle &c1, const Circle &c2, int sign1){
                                                                 tree[M].x1 = tree[M].x2 = tree[M].x;
                                                                 tree[M].y1 = tree[M].y2 = tree[M].y
// sign1 = 1 for outer tang, -1 for inter tang
vector<Line> ret;
                                                                 tree[M].L = build_tree(L, M-1, d+1);
 if (norm(c1.o - c2.o) < eps) return ret;</pre>
                                                                 if (tree[M].L) {
11f'd = abs(c1.o - c2.o);
                                                                  tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
                                                                  tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
PTF v = (c2.0 - c1.0) / d;
                                                                  tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
llf c = (c1.r - sign1 * c2.r) / d;
if (c * c > 1) return ret;
llf h = sqrt(max<llf>(0, 1 - c * c));
for (int sign2 : {1, -1}) {
                                                                 tree[M].R = build_tree(M+1, R, d+1);
 PTF n = c * v + sign2 * h * rot90(v);
                                                                 if (tree[M].R) {
 PTF p1 = c1.o + n * c1.r;
                                                                  tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
 PTF p2 = c2.o + n * (c2.r * sign1);
                                                                  tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
 if (norm(p2 - p1) < eps)
  p2 = p1 + rot90(c2.o - c1.o);
 ret.push_back({p1, p2});
                                                                 return tree+M;
return ret;
                                                                int touch(Node* r, int x, int y, LL d2){
                                                                 LL dis = sqrt(d2)+1;
6.16 Minimum Covering Circle
                                                                 if (x<r->x1-dis || x>r->x2+dis ||
                                                                   y<r->y1-dis || y>r->y2+dis)
template<typename P>
Circle getCircum(const P &a, const P &b, const P &c){
                                                                  return 0;
Real a1 = a.x-b.x, b1 = a.y-b.y;
                                                                 return 1;
Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
                                                                void nearest(Node* r,int x,int y,int &mID,LL &md2) {
Real a2 = a.x-c.x, b2 = a.y-c.y;
Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
                                                                 if (!r || !touch(r, x, y, md2)) return;
Circle cc;
                                                                 LL d2 = dis2(r->x, r->y, x, y);
cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
                                                                 if (d2 < md2 \mid \mid (d2 == md2 \&\& mID < r->id)) {
cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
                                                                  mID = r -> id;
                                                                  md2 = d2:
cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
return cc;
                                                                 }
}
                                                                 // search order depends on split dim
                                                                 if ((r->f == 0 \&\& x < r->x) ||
template<typename P>
                                                                   (r->f == 1 \&\& y < r->y)) {
Circle MinCircleCover(const vector<P>& pts){
                                                                  nearest(r\rightarrow L, x, y, mID, md2);
 random_shuffle(pts.begin(), pts.end());
                                                                  nearest(r->R, x, y, mID, md2);
Circle c = { pts[0], 0 };
                                                                 } else {
for(int i=0;i<(int)pts.size();i++){</pre>
                                                                  nearest(r->R, x, y, mID, md2);
nearest(r->L, x, y, mID, md2);
 if (dist(pts[i], c.o) <= c.r) continue;</pre>
 c = { pts[i], 0 };
 for (int j = 0; j < i; j++) {
   if(dist(pts[j], c.o) <= c.r) continue;</pre>
                                                                int query(int x, int y) {
   c.o = (pts[i] + pts[j]) / 2;
```

int id = 1029384756;

```
LL d2 = 102938475612345678LL;
                                                                       if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
  nearest(root, x, y, id, d2);
                                                                       else{
  return id;
                                                                        sort(eve, eve + E);
} tree;
                                                                        eve[E] = eve[0];
                                                                         for(int j = 0; j < E; ++j){
6.18
       Rotating Sweep Line
                                                                         cnt += eve[j].add;
                                                                         Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
void rotatingSweepLine(pair<int, int> a[], int n) {
 vector<pair<int, int>> 1;
                                                                          double theta = eve[j + 1].ang - eve[j].ang;
                                                                          if (theta < 0) theta += 2. * pi;</pre>
 1.reserve(n * (n - 1) / 2)
 for (int i = 0; i < n; ++i)
                                                                          Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
  for (int j = i + 1; j < n; ++j)
   1.emplace_back(i, j);
 sort(1.begin(), 1.end(), [&a](auto &u, auto &v){
                                                                      }
  1ld udx = a[u.first].first - a[u.second].first;
  1ld udy = a[u.first].second - a[u.second].second;
                                                                   };
  1ld vdx = a[v.first].first - a[v.second].first;
1ld vdy = a[v.first].second - a[v.second].second;
                                                                    7
                                                                         Stringology
  if (udx == 0 or vdx == 0) return not udx == 0;
                                                                    7.1 Suffix Array
  int s = sgn(udx * vdx);
  return udy * vdx * s < vdy * udx * s;
                                                                   namespace sfx {
                                                                   bool _t[maxn * 2];
 });
 vector<int> idx(n), p(n);
                                                                   int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], _c[maxn * 2];
int x[maxn], _p[maxn], _q[maxn * 2];
// sa[i]: sa[i]-th suffix is the
 iota(idx.begin(), idx.end(), 0);
 sort(idx.begin(), idx.end(), [&a](int i, int j){
 return a[i] < a[j]; });
for (int i = 0; i < n; ++i) p[idx[i]] = i;
                                                                    // i-th lexigraphically smallest suffix.
 for (auto [i, j]: 1) {
                                                                    // hi[i]: longest common prefix
  // do here
                                                                    // of suffix sa[i] and suffix sa[i - 1].
  swap(p[i], p[j]);
                                                                    void pre(int *a, int *c, int n, int z) {
                                                                    memset(a, 0, sizeof(int) * n);
  idx[p[i]] = i, idx[p[j]] = j;
                                                                     memcpy(x, c, sizeof(int) * z);
                                                                    void induce(int *a,int *c,int *s,bool *t,int n,int z){
                                                                    memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
if (a[i] && !t[a[i] - 1])
6.19 Circle Cover
const int N = 1021;
struct CircleCover {
 int C;
                                                                       a[x[s[a[i] - 1]]++] = a[i] - 1;
                                                                     memcpy(x, c, sizeof(int) * z);
for (int i = n - 1; i \ge 0; --i)
 Cir c[N]
 bool g[N][N], overlap[N][N];
 // Area[i] : area covered by at least i circles
                                                                      if (a[i] && t[a[i] - 1])
 double Area[ N ];
                                                                       a[--x[s[a[i] - 1]]] = a[i] - 1;
 void init(int _C){ C = _C;}
 struct Teve {
                                                                    void sais(int *s, int *a, int *p, int *q,
  PTF p; double ang; int add;
Teve() {}
                                                                     bool *t, int *c, int n, int z) {
bool uniq = t[n - 1] = true;
  Teve(PTF _a, double _b, int _c):p(_a), ang(_b), add(
                                                                     int nn=0, nmxz=-1, *nsa = a+n, *ns=s+n, last=-1;
                                                                     memset(c, 0, sizeof(int) * z);
     _c){}
  bool operator<(const Teve &a)const
                                                                     for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
                                                                     for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
  {return ang < a.ang;}</pre>
                                                                     if (uniq) +
 }eve[N * 2];
 // strict: x = 0, otherwise x = -1
bool disjuct(Cir &a, Cir &b, int x)
                                                                      for (int i = 0; i < n; ++i) a[--c[s[i]]] = i;
                                                                      return:
 {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
bool contain(Cir &a, Cir &b, int x)
                                                                     for (int i = n - 2; i \ge 0; --i)
 {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
                                                                      t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 bool contain(int i, int j) {
                                                                     pre(a, c, n, z);
                                                                     for (int i = 1; i <= n - 1; ++i)
  /* c[j] is non-strictly in c[i]. */
                                                                      if (t[i] && !t[i - 1])
  return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c
     [j].R) == 0 && i < j)) && contain(c[i], c[j], -1);
                                                                       a[--x[s[i]]] = p[q[i] = nn++] = i;
                                                                     induce(a, c, s, t, n, z);
for (int i = 0; i < n; ++i)
 void solve(){
  fill_n(Area, C + 2, 0);
                                                                      if (a[i] && t[a[i]] && !t[a[i] - 1]) {
  for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)
                                                                      bool neq = last < 0 || \</pre>
                                                                       memcmp(s + a[i], s + last,
(p[q[a[i]] + 1] - a[i]) * sizeof(int));
    overlap[i][j] = contain(i, j);
  for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)
g[i][j] = !(overlap[i][j] || overlap[j][i] ||</pre>
                                                                      ns[q[last = a[i]]] = nmxz += neq;
                                                                     sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
       disjuct(c[i], c[j], -1));
                                                                     pre(a, c, n, z);
  for(int i = 0; i < C; ++i){</pre>
                                                                     for (int i = nn - 1; i >= 0; --i
   int E = 0, cnt = 1;
                                                                      a[--x[s[p[nsa[i]]]] = p[nsa[i]];
   for(int j = 0; j < C; ++j)
                                                                     induce(a, c, s, t, n, z);
    if(j != i && overlap[j][i])
      ++cnt;
                                                                    void build(const string &s) {
   for(int j = 0; j < C; ++j)
                                                                     const int n = int(s.size());
    if(i != j && g[i][j]) {
  auto IP = intersectPoint(c[i], c[j]);
                                                                     for (int i = 0; i < n; ++i) _s[i] = s[i];
                                                                     _s[n] = 0; // s shouldn't contain 0
                                                                     sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
for(int i = 0; i < n; ++i) rev[sa[i] = sa[i+1]] = i;</pre>
      PTF aa = IP[0], bb = IP[1];
      llf A = arg(aa-c[i].0), B = arg(bb-c[i].0);
      eve[E++] = Teve(bb,B,1), eve[E++] = Teve(aa,A,-1);
                                                                     int ind = hi[0] = 0;
      if(B > A) ++cnt;
                                                                     for (int i = 0; i < n; ++i) {
```

**if** (!rev[i]) {

```
ind = 0;
   continue;
                                                               7.4 Manacher
 while (i + ind < n && \</pre>
                                                              int z[maxn]:
   s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
                                                              int manacher(const string& s) {
  string t = ".";
 hi[rev[i]] = ind ? ind-- : 0;
                                                                for(char c: s) t += c, t += '.';
                                                                int 1 = 0, r = 0, ans = 0;
                                                                for (int i = 1; i < t.length(); ++i) {</pre>
7.2 Suffix Automaton
                                                                z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
while (i - z[i] >= 0 && i + z[i] < t.length()) {
struct SuffixAutomaton {
struct node
                                                                  if(t[i - z[i]] == t[i + z[i]]) ++z[i];
  int ch[K], len, fail, cnt, indeg;
                                                                  else break;
  node(int L = 0) : ch{}, len(L), fail(0), cnt(0),
    indeg(0) \{ \}
                                                                 if (i + z[i] > r) r = i + z[i], l = i;
 } st[N];
 int root, last, tot;
                                                                for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
void extend(int c) {
                                                                return ans:
 int cur = ++tot;
                                                              }
  st[cur] = node(st[last].len + 1);
  while (last && !st[last].ch[c]) {
                                                               7.5 Lexico Smallest Rotation
    st[last].ch[c] = cur;
                                                              string mcp(string s) {
    last = st[last].fail;
                                                               int n = s.length();
                                                                s += s; int i = 0, j = 1;
  if (!last) {
                                                                while (i < n && j < n) {</pre>
    st[cur].fail = root;
                                                                 int k = 0:
  } else {
                                                                 while (k < n \&\& s[i + k] == s[j + k]) k++;
    int q = st[last].ch[c];
                                                                 ((s[i+k] \leftarrow s[j+k])?j:i) += k+1;
    if (st[q].len == st[last].len + 1) {
                                                                 j += (i == j);
      st[cur].fail = q;
    } else {
                                                                return s.substr(i < n ? i : j, n);</pre>
      int clone = ++tot;
                                                              }
      st[clone] = st[q];
      st[clone].len = st[last].len + 1;
                                                               7.6 BWT
      st[st[cur].fail = st[q].fail = clone].cnt = 0;
                                                               struct BurrowsWheeler{
      while (last && st[last].ch[c] == q) {
                                                               #define SIGMA 26
        st[last].ch[c] = clone;
                                                               #define BASE 'a'
        last = st[last].fail;
                                                                vector<int> v[ SIGMA ];
      }
                                                                void BWT(char* ori, char* res){
   }
                                                                // make ori -> ori + ori
                                                                 // then build suffix array
 st[last = cur].cnt += 1;
                                                                void iBWT(char* ori, char* res){
void init(const char* s) {
                                                                 for( int i = 0 ; i < SIGMA ; i ++ )</pre>
 root = last = tot = 1;
                                                                  v[ i ].clear();
  st[root] = node(0);
                                                                 int len = strlen( ori );
 for (char c; c = *s; ++s) extend(c - 'a');
                                                                 for( int i = 0 ; i < len ; i ++ )</pre>
                                                                  v[ ori[i] - BASE ].push_back( i );
int q[N];
                                                                 vector<int> a;
void dp() {
                                                                 for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
 for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
                                                                  for( auto j : v[ i ] ){
  a.push_back( j );
  int head = 0, tail = 0;
                                                                   ori[ ptr ++ ] = BASE + i;
 for (int i = 1; i <= tot; i++)
   if (st[i].indeg == 0) q[tail++] = i;
                                                                 for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
 while (head != tail) {
    int now = q[head++];
if (int f = st[now].fail) {
                                                                  ptr = a[ ptr ];
      st[f].cnt += st[now].cnt;
                                                                 res[ len ] = 0;
      if (--st[f].indeg == 0) q[tail++] = f;
    }
                                                              } bwt;
 }
                                                               7.7 Palindromic Tree
int run(const char* s) {
                                                              struct palindromic_tree{
 int now = root;
 for (char c; c = *s; ++s) {
   if (!st[now].ch[c -= 'a']) return 0;
                                                                struct node{
                                                                 int next[26],f,len;
   now = st[now].ch[c];
                                                                 int cnt, num, st, ed;
                                                                 node(int 1=0):f(0),len(1),cnt(0),num(0) {
 return st[now].cnt;
                                                                 memset(next, 0, sizeof(next)); }
} SAM;
                                                                vector<node> st;
                                                                vector<char> s;
7.3 Z value
                                                                int last,n;
vector<int> Zalgo(const string &s) {
                                                                void init(){
vector<int> z(s.size(), s.size());
                                                                 st.clear();s.clear();last=1; n=0;
for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
                                                                 st.push_back(0);st.push_back(-1);
 int j = clamp(r - i, 0, z[i - 1]);
                                                                 st[0].f=1;s.push_back(-1); }
                                                                int getFail(int x){
 for (; i + j < z[0] \text{ and } s[i + j] == s[j]; ++j);
  if (i + (z[i] = j) > r) r = i + z[1 = i];
                                                                 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
                                                                 return x;}
```

return z;

```
void add(int c){
  s.push_back(c-='a'); ++n;
  int cur=getFail(last);
  if(!st[cur].next[c]){
   int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
   st[now].num=st[st[now].f].num+1;
  last=st[cur].next[c];
  ++st[last].cnt;}
 void dpcnt() {
  for (int i=st.size()-1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
int size(){ return st.size()-2;}
} pt:
int main() {
string s; cin >> s; pt.init();
for (int i=0; i<SZ(s); i++) {</pre>
  int prvsz = pt.size(); pt.add(s[i]);
  if (prvsz != pt.size()) {
  int r = i, 1 = r - pt.st[pt.last].len + 1;
   // pal @ [1,r]: s.substr(1, r-1+1)
 }
}
return 0;
```

#### Misc 8

#### 8.1 Theorems

### 8.1.1 Sherman-Morrison formula

$$(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1+v^{\mathsf{T}}A^{-1}u}$$

#### 8.1.2 Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .

### 8.1.3 Tutte's Matrix

Let D be a n imes n matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}.$   $\frac{rank(D)}{2}$  is the maximum matching on G.

# 8.1.4 Cayley's Formula

- Given a degree sequence  $d_1, d_2, \dots, d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1, 2, \ldots, k$  belong to different components. Then  $T_{n,k} =$  $kn^{n-k-1}$ .

#### 8.1.5 Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+$  $\ldots + d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all 1 < k < n.

#### 8.1.6 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

#### 8.1.7 Euler's planar graph formula

$$V - E + F = C + 1, E \le 3V - 6$$
(?)

### 8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have  ${\cal A}$ #{lattice points in the interior}  $+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$ 

#### 8.1.9 Matroid Intersection

Given matroids  $M_1=(G,I_1), M_2=(G,I_2)$ , find maximum  $S\in I_1\cap I_2$ . For each iteration, build the directed graph and find a shortest path from s to t.

```
• s \rightarrow x : S \sqcup \{x\} \in I_1
• x \rightarrow t : S \sqcup \{x\} \in I_2
```

- $y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1$  (y is in the unique circuit of  $S \sqcup \{x\}$ )
- $x \to y : S \setminus \{y\} \sqcup \{x\} \in I_2$  (y is in the unique circuit of  $S \sqcup \{x\}$ )

Alternate the path, and |S| will increase by 1. Let  $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|$ . In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to  $x\in S$  and  $x\notin S$ , resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1

# Bitset LCS

```
scanf("%d%d", &n, &m), u = n / 64 + 1;
for (int i = 1, c; i <= n; i++)
 scanf("%d", &c), p[c].set(i);
for (int i = 1, c; i <= m; i++)
 scanf("%d", &c), (g = f) |= p[c];
 f.shift(), f.set(0);
 ((f = g - f) ^= g) &= g;
printf("%d\n", f.count());
```

```
8.3 Convex 1D/1D DP
struct segment {
int i, 1, r;
 segment() {}
 segment(int a, int b, int c): i(a), l(b), r(c) {}
inline lld f(int l, int r){return dp[l] + w(l+1, r);}
void solve() {
 dp[0] = 0;
 deque<segment> dq; dq.push_back(segment(0, 1, n));
 for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i);
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().l = i + 1
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
   f(i, dq.back().1) < f(dq.back().i, dq.back().1))
    dq.pop_back();
  if (dq.size()) {
   int d = 1 << 20, c = dq.back().1;</pre>
   while (d >>= 1) if (c + d <= dq.back().r)</pre>
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
  if (seg.1 <= n) dq.push_back(seg);</pre>
}
```

# ConvexHull Optimization

```
struct L {
 mutable int64_t a, b, p;
 bool operator<(const L &r) const { return a < r.a; }</pre>
 bool operator<(int64_t x) const { return p < x; }</pre>
};
struct DynamicHull : multiset<L, less<>> {
 static const int64_t kInf = 1e18;
 bool Isect(iterator x, iterator y)
  auto Div = [](int64_t a, int64_t b) {
    return a / b - ((a ^ b) < 0 && a % b); }
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(int64_t a, int64_t b) {
  auto z = insert(\{a, b, 0\}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
  while ((y = x) != begin() \& (--x)->p >= y->p)
   Isect(x, erase(y));
 int64_t Query(int64_t x) {
 auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
};
```

# 8.5 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
int s = 0:
for (int i = 2; i <= n; i++)
 s = (s + m) \% i;
 return s;
// died at kth
int kth(int n, int m, int k){
if (m == 1) return n-1;
for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
return k;
8.6 Cactus Matching
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
dfn[u]=low[u]=++dfs_idx;
for(int i=0;i<(int)init_g[u].size();i++){</pre>
  int v=init_g[u][i];
  if(v==par[u]) continue;
  if(!dfn[v]){
  par[v]=u;
   tarjan(v);
   low[u]=min(low[u],low[v]);
   if(dfn[u]<low[v]){</pre>
    g[u].push_back(v);
    g[v].push_back(u);
  }else{
   low[u]=min(low[u],dfn[v]);
   if(dfn[v]<dfn[u]){</pre>
    int temp_v=u;
    bcc_id++;
    while(temp_v!=v){
     g[bcc_id+n].push_back(temp_v);
     g[temp_v].push_back(bcc_id+n);
     temp_v=par[temp_v];
    g[bcc_id+n].push_back(v);
    g[v].push_back(bcc_id+n);
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
int dp[maxn][2],min_dp[2][2],tmp[2][2],tp[2];
void dfs(int u,int fa){
if(u<=n){
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
   dfs(v,u)
   memset(tp,0x8f,sizeof tp);
   if(v<=n){
    tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
    tp[1]=max(
     dp[u][0]+dp[v][0]+1
     dp[u][1]+max(dp[v][0],dp[v][1])
   }else{
    tp[0]=dp[u][0]+dp[v][0];
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
   dp[u][0]=tp[0],dp[u][1]=tp[1];
 }else{
 for(int i=0;i<(int)g[u].size();i++){</pre>
  int v=g[u][i];
   if(v==fa) continue;
  dfs(v,u);
 min_dp[0][0]=0;
 min_dp[1][1]=1;
min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f3f;
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
```

memset(tmp,0x8f,sizeof tmp);

```
tmp[0][0]=max(
    min_dp[0][0]+max(dp[v][0],dp[v][1]),
    min_dp[0][1]+dp[v][0]
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
   tmp[1][0]=max(
    \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
    min_dp[1][1]+dp[v][0]
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
   memcpy(min_dp,tmp,sizeof tmp);
  dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
  dp[u][0]=min_dp[0][0];
int main(){
 int m,a,b;
 scanf("%d%d",&n,&m);
 for(int i=0;i<m;i++){</pre>
  scanf("%d%d",&a,&b);
  init_g[a].push_back(b);
  init_g[b].push_back(a);
 par[1]=-1;
 tarjan(1);
 dfs(1,-1);
 printf("%d\n", max(dp[1][0], dp[1][1]));
 return 0:
8.7 Tree Knapsack
int dp[N][K]; PII obj[N];
vector<int> G[N];
void dfs(int u,
                int mx){
 for(int s: G[u]) {
  if(mx < obj[s].first) continue;</pre>
  for(int i=0;i<=mx-obj[s].FF;i++)</pre>
   dp[s][i] = dp[u][i];
  dfs(s, mx - obj[s].first);
  for(int i=obj[s].FF;i<=mx;i++)</pre>
   dp[u][i] = max(dp[u][i],
    dp[s][i - obj[s].FF] + obj[s].SS);
}
8.8 N Queens Problem
vector< int > solve( int n ) {
 // no solution when n=2, 3
 vector< int > ret:
 if ( n % 6 == 2 ) {
for ( int i = 2 ; i <= n ; i += 2 )
   ret.push_back( i )
  ret.push_back( 3 ); ret.push_back( 1 );
  for ( int i = 7 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 5 );
 } else if ( n % 6 == 3 ) {
for ( int i = 4 ; i <= n ; i += 2 )</pre>
   ret.push_back( i )
  ret.push_back( 2 );
  for ( int i = 5 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 1 ); ret.push_back( 3 );
 } else {
  for ( int i = 2 ; i <= n ; i += 2 )
   ret.push_back( i );
  for ( int i = 1 ; i <= n ; i += 2 )
   ret.push_back( i );
 return ret;
8.9 Binary Search On Fraction
struct Q {
 11 p, q;
Q go(Q b, 11 d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 \le p,q \le N
```

```
Q frac_bs(ll N) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)
    if (Q mid = hi.go(lo, len + step);
        mid.p > N || mid.q > N || dir ^ pred(mid))
        t++;
    else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  }
  return dir ? hi : lo;
}
```