Contents

1 B						
	asic	1	5.19	Pollard Rho	14	
1.1	vimrc	1	5.20	Berlekamp Massey	14	
1.2	Debug Macro	1		Charateristic Poly-		
1.3	Increase Stack	1		nomial	14	
1.4	Pragma Optimization	i	5.22	Polynomial Opera-		
1.5	IO Optimization	1		tions	15	
		-	5.23	Simplex	15	
2 D	ata Structure	1		Simplex Construction	16	
2.1	Dark Magic	1				
2.2	Link-Cut Tree	2	6 6	eometry	16	
2.3	LiChao Segment Tree	2	6.1	Basic Geometry	16	
2.4	Treap	2	6.2	2D Convex Hull	16	
2.5	Linear Basis	3	6.3	2D Farthest Pair	16	
2.6	Binary Search On	5	6.4	MinMax Enclosing		
2.0	Segtree	3		Rect	16	
			6.5	Minkowski Sum	17	
3 G	raph	3	6.6	Segment Intersection	17	
3.1	2-SAT (SCC)	3	6.7	Half Plane Intersec-		
3.2	BCC	3		tion	17	
3.3	Round Square Tree	4	6.8	SegmentDist		
3.4	Edge TCC	4		(Sausage)	17	
3.5	Centroid Decom-	•	6.9	Rotating Sweep Line	17	
5.5	position	4	6.10	Point In Simple		
3.6	DMST	4		Polygon	17	
3.7	Dominator Tree	5	6.11	Point In Hull (Fast)	17	
3.8		5	6.12	` '	• •	
	Edge Coloring	5	0.12	Tangent of Points To Hull	18	
3.9	Lowbit Decompo-	5	6.13	Circle Class & Inter-	10	
710	sition	_	0.15	section	18	
	Manhattan MST	6 6	6.14	Circle Common		
3.11	MaxClique	О		Tangent	18	
3.12		6	6.15	Line-Circle Inter-		
717	Cycle	O		section	18	
3.13	Mo's Algorithm on	7	6.16	Poly-Circle Inter-		
71/	Tree	-		section	18	
	Tree Hashing	7	6.17	Minimum Covering		
3.15	Virtual Tree	7		Circle	18	
4 M	atching & Flow	7		Circle Union	18	
4.1	HopcroftKarp	7	6.19	Polygon Union	19	
4.2	Dijkstra Cost Flow	7		3D Convex Hull	19	
			6.21	Delaunay	19	
4.3	Dinic	7	6.22	kd Tree (Nearest		
4.4	Flow Models	8		Point)	20	
	C					
4.5	General Graph	0	6.23	kd Closest Pair (3D		
4.5	Matching	8	6.23	kd Closest Pair (3D ver.)	20	
4.5 4.6	Matching	9		ver.)		
4.5	Matching	9	6.24	ver.) Simulated Annealing	20	
4.5 4.6 4.7 4.8	Matching	9	6.24	ver.)		
4.5 4.6 4.7	Matching	9 9 9	6.24	ver.)	20	
4.5 4.6 4.7 4.8 4.9	Matching	9	6.24 7 St	ver.)	20 20	
4.5 4.6 4.7 4.8 4.9	Matching	9 9 9	6.24 7 St 7.1 7.2	ver.)	20 20 20 21	
4.5 4.6 4.7 4.8 4.9 4.10	Matching	9 9 9 9	6.24 7 St 7.1 7.2 7.3	ver.)	20 20 20 21 21	
4.5 4.6 4.7 4.8 4.9 4.10 4.11	Matching	9 9 9 9 10 10	6.24 7 St 7.1 7.2 7.3 7.4	ver.)	20 20 20 21 21 21	
4.5 4.6 4.7 4.8 4.9 4.10	Matching	9 9 9 9 10 10	6.24 7 St 7.1 7.2 7.3	ver.)	20 20 20 21 21	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1	Matching	9 9 9 9 10 10 11	6.24 7 St 7.1 7.2 7.3 7.4 7.5	ver.)	20 20 20 21 21 21	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M	Matching	9 9 9 9 10 10	6.24 7 St 7.1 7.2 7.3 7.4 7.5	ver.) Simulated Annealing Eringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Ro-	20 20 21 21 21 21	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1	Matching	9 9 9 9 10 10 11	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6	ver.) Simulated Annealing tringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation	20 20 21 21 21 21 21	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3	Matching	9 9 9 9 10 10 11 11 11 12	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6	ver.) Simulated Annealing tringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz	20 20 21 21 21 21 21 22 22	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4	Matching	9 9 9 9 10 10 11 11 11 12 12	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree	20 20 21 21 21 22 22 22 22 22 22	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5	Matching	9 9 9 9 10 10 11 11 11 12 12 12	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M	ver.) Simulated Annealing tringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc	20 20 21 21 21 22 22 22 22 22 22	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6	Matching	9 9 9 9 10 10 11 11 12 12 12 12	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1	ver.) Simulated Annealing tringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems	20 20 21 21 21 22 22 22 22 22 22	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7	Matching	9 9 9 9 10 10 11 11 12 12 12 12 12	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid In-	20 20 20 21 21 21 22 22 22 22 22 22 22	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8	Matching	9 9 9 9 10 10 11 11 12 12 12 12 12 12 12	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2	ver.) Simulated Annealing tringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection	20 20 20 21 21 21 22 22 22 22 22 22 22 22 22 22	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9	Matching	9 9 9 9 10 10 11 11 12 12 12 12 12 12	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS	200 200 210 210 210 210 220 220 220 220	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10	Matching	9 9 9 10 10 11 11 12 12 12 12 12 12 12 12 12	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3 8.4	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS Prefix Substring LCS.	200 200 211 211 212 222 222 222 223 233 233	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11	Matching	9 9 9 9 10 10 11 11 12 12 12 12 12 12	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3 8.4 8.5	ver.) Simulated Annealing tringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS Prefix Substring LCS Convex 1D/1D DP	200 200 210 210 210 210 220 220 220 220	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10	Matching	9 9 9 9 10 10 11 11 12 12 12 12 12 12 12 13	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3 8.4	ver.) Simulated Annealing tringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS Prefix Substring LCS Convex 1D/1D DP ConvexHull Opti-	200 200 210 210 210 220 220 220 220 230 230 230 230 230 23	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12	Matching	9 9 9 9 10 10 11 11 12 12 12 12 12 12 12 13	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3 8.4 8.5 8.6	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS Prefix Substring LCS Convex 1D/1D DP ConvexHull Optimization	200 200 201 211 212 222 222 222 223 233 233 233 23	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12	Matching	9 9 9 9 10 10 11 11 12 12 12 12 12 12 12 13 13	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3 8.4 8.5 8.6 8.7	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS Prefix Substring LCS Convex 1D/1D DP ConvexHull Optimization Josephus Problem	200 200 201 211 212 222 222 222 223 233 233 233 23	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14	Matching	9 9 9 10 10 11 11 12 12 12 12 12 12 12 13 13 13	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3 8.4 8.5 8.6	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS Prefix Substring LCS Convex ID/ID DP ConvexHull Optimization Josephus Problem Tree Knapsack	200 200 201 211 212 222 222 222 223 233 233 244	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15	Matching	9 9 9 10 10 11 11 12 12 12 12 12 12 12 13 13 13 13	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3 8.4 8.5 8.6 8.7	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS Prefix Substring LCS Convex 1D/1D DP ConvexHull Optimization Josephus Problem	200 200 201 211 212 222 222 222 223 233 233 233 23	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16	Matching	9 9 9 10 10 11 11 12 12 12 12 12 12 12 13 13 13	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3 8.4 8.5 8.6 8.7 8.8	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS Prefix Substring LCS Convex 1D/1D DP ConvexHull Optimization Josephus Problem Tree Knapsack N Queens Problem	200 200 201 211 212 222 222 222 223 233 233 244	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15	Matching	9 9 9 10 10 11 11 12 12 12 12 12 12 12 13 13 13 14	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3 8.4 8.5 8.6 8.7 8.8 8.9	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS Prefix Substring LCS Convex 1D/1D DP ConvexHull Optimization Josephus Problem Tree Knapsack N Queens Problem Stable Marriage	20 20 21 21 21 22 22 22 22 23 23 23 23 24 24	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17	Matching	9 9 9 10 10 11 11 12 12 12 12 12 12 12 13 13 13 13 14	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3 8.4 8.5 8.6 8.7 8.8 8.9 8.10	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS Prefix Substring LCS Convex ID/ID DP ConvexHull Optimization Josephus Problem Tree Knapsack N Queens Problem	20 20 21 21 21 22 22 22 22 23 23 23 23 24 24	
4.5 4.6 4.7 4.8 4.9 4.10 4.11 5 M 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17	Matching	9 9 9 10 10 11 11 12 12 12 12 12 12 12 13 13 13 14	6.24 7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 7.9 8 M 8.1 8.2 8.3 8.4 8.5 8.6 8.7 8.8 8.9 8.10	ver.) Simulated Annealing ringology Hash Suffix Array Ex SAM Z value Manacher Lexico Smallest Rotation Main Lorentz BWT Palindromic Tree isc Theorems Weight Matroid Intersection Bitset LCS Prefix Substring LCS Convex 1D/1D DP ConvexHull Optimization Josephus Problem Tree Knapsack N Queens Problem Stable Marriage Binary Search On	20 20 21 21 21 22 22 22 22 23 23 23 24 24 24	

1 Basic

1.1 vimrc

```
se is nu ru et tgc sc hls cin cino+=j1 sw=2 sts=2 bs=2
mouse=a "encoding=utf-8 ls=2
```

1.2 Debug Macro [b78d75]

```
#ifdef CKISEKI
#define safe cerr<<__PRETTY_FUNCTION__<<" line "<</pre>
      _LINE__<<" safe\n"
#define debug(a...) debug_(#a, a)
#define orange(a...) orange_(#a, a)
template <typename ...T>
void debug_(const char *s, T ...a) {
  cerr << "\e[1;32m(" << s << ") = (";</pre>
  int cnt = sizeof...(T);
  (..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
template <typename I>
void orange_(const char *s, I L, I R) {
   cerr << "\e[1;32m[ " << s << " ] = [ "</pre>
  for (int f = 0; L != R; ++L)
    cerr << (f++ ? ", " : "") << *L;
  cerr << " ]\e[0m\n";</pre>
}
#else
#define safe ((void)0)
#define debug(...) safe
#define orange(...) safe
#endif
```

1.3 Increase Stack

```
const int size = 256 << 20;
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));</pre>
```

1.4 Pragma Optimization [f63b0a]

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8000)
```

1.5 IO Optimization [c9494b]

```
static inline int gc() {
  constexpr int B = 1<<20; static char buf[B], *p, *q;
  if (p == q) q = (p = buf) + fread(buf, 1, B, stdin);
  return q == buf ? EOF : *p++;
}</pre>
```

2 Data Structure

2.1 Dark Magic [095f25]

2.2 Link-Cut Tree [7ce2b4]

```
template <typename Val, typename SVal> class LCT {
struct node {
 int pa, ch[2];
 bool rev;
 Val v, prod, rprod;
 SVal sv, sub, vir;
 node(): pa{0}, ch{0, 0}, rev{false}, v{}, prod{},
    rprod{}, sv{}, sub{}, vir{} {};
#define cur o[u]
#define lc cur.ch[0]
#define rc cur.ch[1]
vector<node> o;
bool is_root(int u) const {
 return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u;
bool is_rch(int u) const {
 return o[cur.pa].ch[1] == u && !is_root(u);
void down(int u) {
 if (not cur.rev) return;
 if (lc) set_rev(lc);
 if (rc) set_rev(rc);
 cur.rev = false;
void up(int u) {
 cur.prod = o[lc].prod * cur.v * o[rc].prod;
 cur.rprod = o[rc].rprod * cur.v * o[lc].rprod;
 cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
void set_rev(int u) {
 swap(lc, rc);
 swap(cur.prod, cur.rprod);
 cur.rev ^= 1;
void rotate(int u) {
 int f=cur.pa, g=o[f].pa, l=is_rch(u);
if (cur.ch[l ^ 1]) o[cur.ch[l ^ 1]].pa = f;
 if (not is_root(f)) o[g].ch[is_rch(f)] = u;
 o[f].ch[l] = cur.ch[l ^ 1];
 cur.ch[l ^ 1] = f;
 cur.pa = g, o[f].pa = u;
 up(f);
void splay(int u) {
 vector<int> stk = {u};
 while (not is_root(stk.back()))
  stk.push_back(o[stk.back()].pa);
 while (not stk.empty()) {
  down(stk.back());
  stk.pop_back();
 for (int f = cur.pa; not is_root(u); f = cur.pa) {
  if(!is_root(f))rotate(is_rch(u)==is_rch(f)?f:u);
  rotate(u);
 }
 up(u);
void access(int x) {
 for (int u = x, last = 0; u; u = cur.pa) {
  splay(u);
  cur.vir = cur.vir + o[rc].sub - o[last].sub;
  rc = last; up(last = u);
 }
 splay(x);
int find_root(int u) {
 int la = 0;
 for (access(u); u; u = lc) down(la = u);
 return la;
void split(int x, int y) {change_root(x);access(y);}
void change_root(int u) { access(u); set_rev(u); }
public:
LCT(int n = 0) : o(n + 1) {}
int add(const Val &v = {}) {
 o.push_back(v);
 return int(o.size()) - 2;
int add(Val &&v) {
```

```
o.emplace_back(move(v));
  return int(o.size()) - 2;
 void set_val(int u, const Val &v) {
 splay(++u); cur.v = v; up(u);
 void set_sval(int u, const SVal &v) {
 splay(++u); cur.sv = v; up(u);
 Val query(int x, int y) {
 split(++x, ++y); return o[y].prod;
 SVal subtree(int p, int u) {
 change_root(++p); access(++u);
  return cur.vir + cur.sv;
 bool connected(int u, int v) {
 return find_root(++u) == find_root(++v); }
 void link(int x, int y) {
 change_root(++x); access(++y);
 o[y].vir = o[y].vir + o[x].sub;
 up(o[x].pa = y);
 void cut(int x, int y) {
  split(++x,
             ++v);
 o[y].ch[0] = o[x].pa = 0; up(y);
#undef cur
#undef lc
#undef rc
```

2.3 LiChao Segment Tree [b9c827]

```
struct L {
 int m, k, id;
 L(): id(-1) {}
 L(int a, int b, int c) : m(a), k(b), id(c) {}
 int at(int x) { return m * x + k; }
class LiChao {
private:
 int n; vector<L> nodes;
 static int lc(int x) { return 2 * x + 1; }
 static int rc(int x) { return 2 * x + 2; }
 void insert(int l, int r, int id, L ln) {
  int m = (l + r) >> 1;
  if (nodes[id].id == -1)
   return nodes[id] = ln, void();
  bool atLeft = nodes[id].at(l) < ln.at(l);</pre>
  if (nodes[id].at(m) < ln.at(m))</pre>
   atLeft ^= 1, swap(nodes[id], ln);
  if (r - l == 1) return;
  if (atLeft) insert(l, m, lc(id), ln);
  else insert(m, r, rc(id), ln);
 int query(int l, int r, int id, int x) {
  int m = (l + r) >> 1, ret = 0;
  if (nodes[id].id != -1) ret = nodes[id].at(x);
  if (r - l == 1) return ret;
  if (x < m) return max(ret, query(l, m, lc(id), x));</pre>
  return max(ret, query(m, r, rc(id), x));
public:
LiChao(int n_{-}) : n(n_{-}), nodes(n * 4) {}
 void insert(L ln) { insert(0, n, 0, ln); }
 int query(int x) { return query(0, n, 0, x); }
};
```

2.4 Treap [ae576c]

```
__gnu_cxx::sfmt19937 rnd(7122);
namespace Treap {
struct node {
  int size, pri; node *lc, *rc, *pa;
  node() : size(1), pri(rnd()), lc(0), rc(0), pa(0) {}
  void pull() {
    size = 1; pa = 0;
    if (lc) { size += lc->size; lc->pa = this; }
    if (rc) { size += rc->size; rc->pa = this; }
  }
};
int SZ(node *x) { return x ? x->size : 0; }
node *merge(node *L, node *R) {
```

```
if (not L or not R) return L ? L : R;
if (L->pri > R->pri)
  return L->rc = merge(L->rc, R), L->pull(), L;
else
 return R->lc = merge(L, R->lc), R->pull(), R;
void splitBySize(node *o, int k, node *&L, node *&R) {
if (not 0) L = R = 0;
else if (int s = SZ(o->lc) + 1; s <= k)
 L=o, splitBySize(o->rc, k-s, L->rc, R), L->pull();
else
 R=o, splitBySize(o->lc, k, L, R->lc), R->pull();
} // SZ(L) == k
int getRank(node *o) { // 1-base
int r = SZ(o->lc) + 1;
for (; o->pa; o = o->pa)
 if (o->pa->rc == o) r += SZ(o->pa->lc) + 1;
return r:
} // namespace Treap
```

2.5 Linear Basis [138d5d]

```
template <int BITS, typename S = int> struct Basis {
  static constexpr S MIN = numeric_limits<S>::min();
 array<pair<llu, S>, BITS> b;
 Basis() { b.fill({0, MIN}); }
 void add(llu x, S p) {
  for (int i = BITS-1; i>=0; i--) if (x >> i & 1) {
   if (b[i].first == 0) return b[i]={x, p}, void();
   if (b[i].second < p)</pre>
    swap(b[i].first, x), swap(b[i].second, p);
   x ^= b[i].first;
 }
 optional<llu> query_kth(llu v, llu k) {
  vector<pair<llu, int>> o;
  for (int i = 0; i < BITS; i++)</pre>
   if (b[i].first) o.emplace_back(b[i].first, i);
  if (k >= (1ULL << o.size())) return {};</pre>
  for (int i = int(o.size()) - 1; i >= 0; i--)
   if ((k >> i & 1) ^ (v >> o[i].second & 1))
    v ^= o[i].first;
  return v;
 Basis filter(S l) {
  Basis res = *this;
  for (int i = 0; i < BITS; i++)</pre>
   if (res.b[i].second < l) res.b[i] = {0, MIN};</pre>
  return res;
 }
};
```

2.6 Binary Search On Segtree [29b3cb]

```
// find_first = x \rightarrow minimal x s.t. check([a, x))
// find_last = x \rightarrow maximal x s.t. check( [x, b) )
template <typename C>
int find_first(int l, const C &check) {
if (l >= n) return n + 1;
for (int i = hei; i > 0; i--) propagate(l >> i);
Monoid sum = identity;
do {
 while ((l & 1) == 0) l >>= 1;
 if (check(f(sum, data[l]))) {
  while (l < sz) {</pre>
    propagate(l); l <<= 1;</pre>
    if (auto nxt = f(sum,data[l]); not check(nxt))
     sum = nxt, l++;
  return l + 1 - sz;
 }
  sum = f(sum, data[l++]);
} while ((l & -l) != l);
return n + 1;
template <typename C>
int find_last(int r, const C &check) {
  if (r <= 0) return -1;</pre>
r += sz;
for (int i = hei; i > 0; i--) propagate((r-1) >> i);
Monoid sum = identity;
do {
```

```
r--;
while (r > 1 and (r & 1)) r >>= 1;
if (check(f(data[r], sum))) {
    while (r < sz) {
        propagate(r); r = (r << 1) + 1;
        if (auto nxt = f(data[r], sum); not check(nxt))
            sum = nxt, r--;
        }
    return r - sz;
    }
    sum = f(data[r], sum);
}    while ((r & -r) != r);
    return -1;
}</pre>
```

3 Graph

3.1 2-SAT (SCC) [76434f]

```
class TwoSat { // test @ CSES Giant Pizza
private:
 int n; vector<vector<int>> G, rG, sccs;
 vector<int> ord, idx, vis, res;
 void dfs(int u) {
  vis[u] = true;
  for (int v : G[u]) if (!vis[v]) dfs(v);
  ord.push_back(u);
 void rdfs(int u) {
  vis[u] = false; idx[u] = sccs.size() - 1;
  sccs.back().push_back(u);
  for (int v : rG[u]) if (vis[v]) rdfs(v);
 }
public:
 TwoSat(int n_{-}): n(n_{-}), G(n), rG(n), idx(n), vis(n),
    res(n) {}
 void add_edge(int u, int v) {
 G[u].push_back(v); rG[v].push_back(u);
 void orr(int x, int y) {
  if ((x ^ y) == 1) return;
  add_edge(x ^ 1, y); add_edge(y ^ 1, x);
 bool solve() {
  for (int i = 0; i < n; ++i) if (not vis[i]) dfs(i);</pre>
  reverse(ord.begin(), ord.end());
  for (int u : ord)
   if (vis[u]) sccs.emplace_back(), rdfs(u);
  for (int i = 0; i < n; i += 2)</pre>
   if (idx[i] == idx[i + 1]) return false;
  vector<bool> c(sccs.size());
  for (size_t i = 0; i < sccs.size(); ++i)</pre>
   for (int z : sccs[i])
    res[z] = c[i], c[idx[z ^ 1]] = !c[i];
  return true;
 bool get(int x) { return res[x]; }
 int get_id(int x) { return idx[x]; }
 int count() { return sccs.size(); }
```

3.2 BCC [6ac6db]

```
class BCC {
 int n, ecnt, bcnt;
 vector<vector<pair<int, int>>> g;
 vector<int> dfn, low, bcc, stk;
 vector<bool> ap, bridge;
void dfs(int u, int f) {
  dfn[u] = low[u] = dfn[f] + 1;
  int ch = 0;
  for (auto [v, t] : g[u]) if (bcc[t] == -1) {
   bcc[t] = 0; stk.push_back(t);
   if (dfn[v]) {
    low[u] = min(low[u], dfn[v]);
    continue;
   ++ch, dfs(v, u);
   low[u] = min(low[u], low[v]);
   if (low[v] > dfn[u]) bridge[t] = true;
   if (low[v] < dfn[u]) continue;</pre>
   ap[u] = true;
   while (not stk.empty()) {
```

```
int o = stk.back(); stk.pop_back();
    bcc[o] = bcnt;
    if (o == t) break;
  bcnt += 1;
 ap[u] = ap[u] and (ch != 1 or u != f);
public:
BCC(int n_) : n(n_), ecnt(0), bcnt(0), g(n), dfn(n),
low(n), stk(), ap(n) {}
void add_edge(int u, int v) {
 g[u].emplace_back(v, ecnt);
  g[v].emplace_back(u, ecnt++);
void solve() {
 bridge.assign(ecnt, false); bcc.assign(ecnt, -1);
 for (int i = 0; i < n; ++i) if (!dfn[i]) dfs(i, i);</pre>
 int bcc_id(int x) const { return bcc[x]; }
bool is_ap(int x) const { return ap[x]; }
bool is_bridge(int x) const { return bridge[x]; }
```

3.3 Round Square Tree [528440]

```
struct RST {
 int n; vector<vector<int>> T;
RST(auto &G) : n(G.size()), T(n) {
 vector<int> stk, vis(n), low(n);
auto dfs = [&](auto self, int u, int d) -> void {
  low[u] = vis[u] = d; stk.push_back(u);
   for (int v : G[u]) if (!vis[v]) {
    self(self, v, d + 1);
    if (low[v] == vis[u]) {
     int cnt = T.size(); T.emplace_back();
     for (int x = -1; x != v; stk.pop_back())
      T[cnt].push_back(x = stk.back());
    T[u].push_back(cnt); // T is rooted
    } else low[u] = min(low[u], low[v]);
  } else low[u] = min(low[u], vis[v]);
  };
 for (int u = 0; u < N; u++)
  if (!vis[u]) dfs(dfs, u, 1);
} // T may be forest; after dfs, stk are the roots
}; // test @ 2020 Shanghai K
```

3.4 Edge TCC [5a2668]

```
vector<vector<int>> ETCC(auto &adj) {
const int n = static_cast<int>(adj.size());
vector<int> up(n), low(n), in, out, nx, id;
in = out = nx = id = vector < int > (n, -1);
 int dfc = 0, cnt = 0; Dsu dsu(n);
auto merge = [&](int u, int v) {
 dsu.join(u, v); up[u] += up[v]; };
auto dfs = [&](auto self, int u, int p) -> void {
  in[u] = low[u] = dfc++
  for (int v : adj[u]) if (v != u) {
   if (v == p) { p = -1; continue; }
   if (in[v] == -1) {
   self(self, v, u);
if (nx[v] == -1 && up[v] <= 1) {</pre>
     up[u] += up[v]; low[u] = min(low[u], low[v]);
     continue:
    if (up[v] == 0) v = nx[v];
    if (low[u] > low[v])
     low[u] = low[v], swap(nx[u], v);
  for (; v != -1; v = nx[v]) merge(u, v);
} else if (in[v] < in[u]) {</pre>
   low[u] = min(low[u], in[v]); up[u]++;
   } else {
    for (int &x = nx[u]; x != -1 &&
      in[x] \le in[v] \& in[v] < out[x]; x = nx[x])
     merge(u, x);
    up[u]--;
  }
 }
  out[u] = dfc;
for (int i = 0; i < n; i++)</pre>
 if (in[i] == -1) dfs(dfs, i, -1);
 for (int i = 0; i < n; i++)</pre>
```

```
if (dsu.anc(i) == i) id[i] = cnt++;
vector<vector<int>> comps(cnt);
for (int i = 0; i < n; i++)
  comps[id[dsu.anc(i)]].push_back(i);
return comps;
} // test @ yosupo judge</pre>
```

3.5 Centroid Decomposition [63b2fb]

```
struct Centroid {
 using G = vector<vector<pair<int, int>>>;
 vector<vector<int64_t>> Dist;
 vector<int> Pa, Dep;
 vector<int64_t> Sub, Sub2;
 vector<int> Cnt, Cnt2;
 vector<int> vis, sz, mx, tmp;
 void DfsSz(const G &g, int x) {
  vis[x] = true, sz[x] = 1, mx[x] = 0;
  for (auto [u, w] : g[x]) if (not vis[u]) {
   DfsSz(g, u); sz[x] += sz[u];
   mx[x] = max(mx[x], sz[u]);
  tmp.push_back(x);
 void DfsDist(const G &g, int x, int64_t D = 0) {
  Dist[x].push_back(D); vis[x] = true;
  for (auto [u, w] : g[x])
   if (not vis[u]) DfsDist(g, u, D + w);
 void DfsCen(const G &g, int x, int D = 0, int p = -1)
  tmp.clear(); DfsSz(g, x);
  int M = tmp.size(), C = -1;
  for (int u : tmp) {
   if (max(M - sz[u], mx[u]) * 2 <= M) C = u;</pre>
   vis[u] = false;
  DfsDist(g, C);
  for (int u : tmp) vis[u] = false;
  Pa[C] = p, vis[C] = true, Dep[C] = D;

for (auto [u, w] : g[C])
   if (not vis[u]) DfsCen(g, u, D + 1, C);
 Centroid(int N, G g)
   : Sub(N), Sub2(N), Cnt(N), Cnt2(N), Dist(N), Pa(N),
    Dep(N), \, vis(N), \, sz(N), \, mx(N) \, \left\{ \, \, \mathsf{DfsCen}(g, \, \theta); \, \, \right\}
 void Mark(int v) {
  int x = v, z = -1;
  for (int i = Dep[v]; i >= 0; --i) {
   Sub[x] += Dist[v][i], Cnt[x]++;
   if (z != -1)
    Sub2[z] += Dist[v][i], Cnt2[z]++;
   x = Pa[z = x];
  }
 int64_t Query(int v) {
  int64_t res = 0;
  int x = v, z = -1;
  for (int i = Dep[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Cnt[x] * Dist[v][i];
   if (z != -1)
    res -= Sub2[z] + 1LL * Cnt2[z] * Dist[v][i];
   x = Pa[z = x];
  return res;
};
```

3.6 DMST [0ae901]

```
using D = int64_t;
struct E { int s, t; D w; }; // 0-base
vector<int> dmst(const vector<E> &e, int n, int root) {
    using PQ = pair<min_heap<pair<D, int>>, D>;
    auto push = [](PQ &pq, pair<D, int> v) {
        pq.first.emplace(v.first - pq.second, v.second);
    };
    auto top = [](const PQ &pq) -> pair<D, int> {
        auto r = pq.first.top();
        return {r.first + pq.second, r.second};
    };
    auto join = [&push, &top](PQ &a, PQ &b) {
        if (a.first.size() < b.first.size()) swap(a, b);
        while (!b.first.empty()) {</pre>
```

```
push(a, top(b));
  b.first.pop();
 }
};
vector<PQ> h(n * 2);
for (size_t i = 0; i < e.size(); ++i)</pre>
push(h[e[i].t], {e[i].w, i});
vector<int> a(n*2), v(n*2, -1), pa(n*2, -1), r(n*2);
iota(a.begin(), a.end(), 0);
auto o = [\&](int x) \{ int y;
 for (y = x; a[y] != y; y = a[y]);
 for (int ox = x; x != y; ox = x)
 x = a[x], a[ox] = y;
 return y;
};
v[root] = n + 1;
int pc = n;
for (int i = 0; i < n; ++i) if (v[i] == -1) {
 for (int p = i; v[p] == -1 || v[p] == i; p = o(e[r[p
   ]].s)) {
  if (v[p] == i) {
   int q = p; p = pc++;
    h[q].second = -h[q].first.top().first;
    join(h[pa[q] = a[q] = p], h[q]);
    while ((q = o(e[r[q]].s)) != p);
  while (!h[p].first.empty() && o(e[top(h[p]).second].
   h[p].first.pop();
  r[p] = top(h[p]).second;
vector<int> ans;
for (int i = pc - 1; i >= 0; i--) if (i != root && v[i
   ] != n) {
 for (int f = e[r[i]].t; f != -1 && v[f] != n; f = pa[
  v[f] = n;
 ans.push_back(r[i]);
return ans; // default minimize, returns edgeid array
```

3.7 Dominator Tree [ea5b7c]

struct Dominator {

```
vector<vector<int>> g, r, rdom; int tk;
vector<int> dfn, rev, fa, sdom, dom, val, rp;
Dominator(int n) : g(n), r(n), rdom(n), tk(0) {
 dfn = rev = fa = sdom = dom =
  val = rp = vector<int>(n, -1); }
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
 rev[dfn[x] = tk] = x;
 fa[tk] = sdom[tk] = val[tk] = tk; tk++;
 for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
  r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
 if (fa[x] == x) return c ? -1 : x;
 if (int p = find(fa[x], 1); p != -1) {
  if (sdom[val[x]] > sdom[val[fa[x]]])
   val[x] = val[fa[x]];
  fa[x] = p;
  return c ? p : val[x];
 } else return c ? fa[x] : val[x];
vector<int> build(int s, int n) {
 // return the father of each node in dominator tree
 dfs(s); // p[i] = -2 \text{ if i is unreachable from s}
 for (int i = tk - 1; i >= 0; --i) {
  for (int u : r[i])
   sdom[i] = min(sdom[i], sdom[find(u)]);
  if (i) rdom[sdom[i]].push_back(i);
  for (int u : rdom[i]) {
   int p = find(u);
   dom[u] = (sdom[p] == i ? i : p);
```

```
if (i) merge(i, rp[i]);
  vector<int> p(n, -2); p[s] = -1;
  for (int i = 1; i < tk; ++i)
if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];</pre>
  for (int i = 1; i < tk; ++i)</pre>
   p[rev[i]] = rev[dom[i]];
  return p;
 } // test @ yosupo judge
}:
3.8 Edge Coloring [029763]
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
 for (int i = 0; i <= N; i++)</pre>
  for (int j = 0; j <= N; j++)
C[i][j] = G[i][j] = 0;</pre>
void solve(vector<pair<int, int>> &E, int N) {
 int X[kN] = {}, a;
 auto update = [&](int u) {
  for (X[u] = 1; C[u][X[u]]; X[u]++);
```

auto color = [&](int u, int v, int c) {

auto flip = [&](int u, int c1, int c2) {

for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {</pre>

if (!C[v][c]) for(a=L.size()-1;a>=0;a--)

color(u, L[a].first, L[a].second);

else if(!C[u][d])**for**(a=L.size()-1;a>=0;a--)

for (; v; v = flip(v, c, d), swap(c, d));
if (C[u][c0]) { a = int(L.size()) - 1;

for(;a>=0;a--)color(u,L[a].first,L[a].second);

while (--a >= 0 && L[a].second != c);

int p = G[u][v];
G[u][v] = G[v][u] = c;
C[u][c] = v, C[v][c] = u;

int p = C[u][c1];

while (!G[u][v0]) {

if (!G[u][v0]) {

} **else** t--;

return p;

return p;

C[u][p] = C[v][p] = 0; if (p) X[u] = X[v] = p;

else update(u), update(v);

swap(C[u][c1], C[u][c2]);

if (!C[u][c1]) X[u] = c1;
if (!C[u][c2]) X[u] = c2;

if (p) G[u][p] = G[p][u] = c2;

L.emplace_back(v, d = X[v]);

else vst[d] = 1, v = C[u][d];

else if (vst[d]) break;

c = color(u, L[a].first, c);

auto [u, v] = E[t]; int v0 = v, c = X[u], c0 = c, d; vector<pair<int, int>> L; int vst[kN] = {};

3.9 Lowbit Decomposition [aa3f57]

```
class LBD {
  int timer, chains;
  vector<vector<int>> G;
  vector<int> tl, tr, chain, head, dep, pa;
  // chains : number of chain
  // tl, tr[u] : subtree interval in the seq. of u
  // head[i] : head of the chain i
  // chian[u] : chain id of the chain u is on
  void predfs(int u, int f) {
    dep[u] = dep[pa[u] = f] + 1;
    for (int v : G[u]) if (v != f) {
        predfs(v, u);
        if (lowbit(chain[u]) < lowbit(chain[v]))
            chain[u] = chain[v];
    }
    if (chain[u] == 0) chain[u] = ++chains;</pre>
```

```
void dfschain(int u, int f) {
 tl[u] = timer++;
  if (head[chain[u]] == -1)
  head[chain[u]] = u;
  for (int v : G[u])
  if (v != f and chain[v] == chain[u])
   dfschain(v, u);
  for (int v : G[u])
   if (v != f and chain[v] != chain[u])
   dfschain(v, u);
  tr[u] = timer;
public:
LBD(int n) : timer(0), chains(0), G(n), tl(n), tr(n),
chain(n), head(n, -1), dep(n), pa(n) \{\} void add_edge(int u, int v) \{
 G[u].push_back(v); G[v].push_back(u);
void decompose() { predfs(0, 0); dfschain(0, 0); }
PII get_subtree(int u) { return {tl[u], tr[u]}; }
vector<PII> get_path(int u, int v) {
 vector<PII> res;
 while (chain[u] != chain[v]) {
  if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
    swap(u, v);
   int s = head[chain[u]];
  res.emplace_back(tl[s], tl[u] + 1);
  u = pa[s];
 if (dep[u] < dep[v]) swap(u, v);</pre>
 res.emplace_back(tl[v], tl[u] + 1);
  return res;
```

3.10 Manhattan MST [df6f59]

```
typedef Point<int> P;
vector<array<int, 3>> manhattanMST(vector<P> ps) {
vi id(sz(ps));
iota(all(id), 0);
vector<array<int, 3>> edges;
rep(k, 0, 4) {
 sort(all(id), [&](int i, int j) {
  return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
 map<int, int> sweep;
 for (int i : id) {
   for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++)) {
    int j = it->second;
    P d = ps[i] - ps[j];
   if (d.y > d.x) break;
    edges.push_back({d.y + d.x, i, j});
  sweep[-ps[i].y] = i;
 for (P &p : ps)
  if (k \& 1) p.x = -p.x;
  else swap(p.x, p.y);
return edges; // [{w, i, j}, ...]
```

3.11 MaxClique [293730]

```
// contain a self loop u to u, than u won't in clique
template <size_t maxn> class MaxClique {
private:
using bits = bitset<maxn>;
bits popped, G[maxn], ans;
size_t deg[maxn], deo[maxn], n;
void sort_by_degree() {
 popped.reset();
  for (size_t i = 0; i < n; ++i)</pre>
  deg[i] = G[i].count();
  for (size_t i = 0; i < n; ++i) {</pre>
   size_t mi = maxn, id = 0;
   for (size_t j = 0; j < n; ++j)</pre>
   if (not popped[j] and deg[j] < mi)</pre>
    mi = deg[id = j];
   popped[deo[i] = id] = 1;
   for (size_t u = G[i]._Find_first(); u < n;</pre>
```

```
u = G[i]._Find_next(u))
     --deg[u];
  }
 void BK(bits R, bits P, bits X) {
  if (R.count() + P.count() <= ans.count()) return;</pre>
  if (not P.count() and not X.count()) {
   if (R.count() > ans.count()) ans = R;
  /* greedily chosse max degree as pivot
  bits cur = P | X; size_t pivot = 0, sz = 0;
  for ( size_t u = cur._Find_first() ;
   u < n ; u = cur.\_Find\_next(u)
    if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
  cur = P & ( ~G[ pivot ] );
  */ // or simply choose first
  bits cur = P & (~G[(P | X)._Find_first()]);
  for (size_t u = cur._Find_first(); u < n;</pre>
    u = cur._Find_next(u)) {
   if (R[u]) continue;
   R[u] = 1;
   BK(R, P \& G[u], X \& G[u]);
   R[u] = P[u] = 0, X[u] = 1;
 }
public:
 void init(size_t n_) {
  n = n_{\cdot}
  for (size_t i = 0; i < n; ++i) G[i].reset();</pre>
  ans.reset();
 void add_edges(int u, bits S) { G[u] = S; }
 void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
 int solve() {
  sort_by_degree(); // or simply iota( deo... )
  for (size_t i = 0; i < n; ++i)</pre>
   deg[i] = G[i].count();
  bits pob, nob = 0; pob.set();
  for (size_t i = n; i < maxn; ++i) pob[i] = 0;</pre>
  for (size_t i = 0; i < n; ++i) {</pre>
   size_t v = deo[i];
   bits tmp;
   tmp[v] = 1;
   BK(tmp, pob \& G[v], nob \& G[v]);
   pob[v] = 0, nob[v] = 1;
  return static_cast<int>(ans.count());
};
```

3.12 Minimum Mean Cycle [e23bc0]

```
// WARNING: TYPE matters
struct Edge { int s, t; llf c; };
llf solve(vector<Edge> &e, int n) {
 // O(VE), returns inf if no cycle, mmc otherwise
 vector<VI> prv(n + 1, VI(n)), prve = prv;
 vector<vector<llf>>> d(n + 1, vector<llf>(n, inf));
 d[0] = vector<llf>(n, 0);
 for (int i = 0; i < n; i++) {</pre>
  for (int j = 0; j < (int)e.size(); j++) {
  auto [s, t, c] = e[j];</pre>
   if (d[i][s] < inf && d[i + 1][t] > d[i][s] + c) {
    d[i + 1][t] = d[i][s] + c;
    prv[i + 1][t] = s; prve[i + 1][t] = j;
 llf mmc = inf; int st = -1;
 for (int i = 0; i < n; i++) {
  llf avg = -inf;
  for (int k = 0; k < n; k++) {
   if (d[n][i] < inf - eps)
    avg = max(avg, (d[n][i] - d[k][i]) / (n - k));
   else avg = inf;
  if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
 if (st == -1) return inf;
 vector<int> vst(n), eid, cycle, rho;
 for (int i = n; !vst[st]; st = prv[i--][st]) {
 vst[st]++; eid.emplace_back(prve[i][st]);
```

```
rho.emplace_back(st);
}
while (vst[st] != 2) {
  int v = rho.back(); rho.pop_back();
  cycle.emplace_back(v); vst[v]++;
}
reverse(all(eid)); eid.resize(cycle.size());
return mmc;
}
```

3.13 Mo's Algorithm on Tree

```
dfs u:
  push u
  iterate subtree
  push u
Let P = LCA(u, v) with St(u)<=St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]</pre>
```

3.14 Tree Hashing [707efa]

```
llu F(llu z) { // xorshift64star from iwiwi
z ^= z >> 12; z ^= z << 25; z ^= z >> 27;
return z * 2685821657736338717LL;
}
llu hsah(int u, int f) {
    llu r = 127; // bigger?
    for (int v : G[u]) if (v != f) r += F( hsah(v, u) );
    return F(r);
} // test @ UOJ 763
```

3.15 Virtual Tree [ad5cf5]

```
vector<pair<int, int>> build(vector<int> vs, int r) {
vector<pair<int, int>> res;
sort(vs.begin(), vs.end(), [](int i, int j) {
  return dfn[i] < dfn[j]; });</pre>
 vector<int> s = {r};
for (int v : vs) if (v != r) {
  if (int o = lca(v, s.back()); o != s.back()) {
  while (s.size() >= 2) {
    if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
    res.emplace_back(s[s.size() - 2], s.back());
    s.pop_back();
   if (s.back() != o) {
    res.emplace_back(o, s.back());
    s.back() = o;
  }
 s.push_back(v);
for (size_t i = 1; i < s.size(); ++i)</pre>
 res.emplace_back(s[i - 1], s[i]);
return res; // (x, y): x->y
```

4 Matching & Flow

4.1 HopcroftKarp [4e7e69]

```
struct HK {
vector<int> l, r, a, p; int ans;
HK(int n, int m, auto \&g) : l(n,-1),r(m,-1),ans(0) {
 for (bool match = true; match; ) {
  match = false; a.assign(n, -1); p = a;
   queue<int> q; int z;
   for (int i = 0; i < n; i++)
    if (l[i] == -1) q.push(a[i] = p[i] = i);
   // bitset<maxn> nvis, t; nvis.set();
   while (!q.empty()) {
    int x = q.front(); q.pop();
   if (l[a[x]] != -1) continue;
    // or use _Find_first and _Find_next here
    for (int y: g[x]) {
     // nvis.reset(y);
     if (r[y] == -1) {
     while (y != -1)
r[y] = x, swap(l[x], y), x = p[x];
      match = true; ans++; break;
     } else if (p[r[y]] == -1)
      q.push(z = r[y]), p[z] = x, a[z] = a[x];
```

```
}
}
};
```

4.2 Dijkstra Cost Flow [06a723]

```
template <typename F, typename C> class MCMF {
 static constexpr F INF_F = numeric_limits<F>::max();
 static constexpr C INF_C = numeric_limits<C>::max();
 struct E {
  int to, r;
  F f; C c;
  E() {}
  E(int a, int b, F x, C y)
   : to(a), r(b), f(x), c(y) {}
 vector<vector<E>> g;
 vector<pair<int, int>> f;
 vector<F> up;
 vector<C> d, h;
 optional<pair<F, C>> step(int S, int T) {
  priority_queue<pair<C, int>> q;
q.emplace(d[S] = 0, S), up[S] = INF_F;
  while (not q.empty()) {
   auto [l, u] = q.top(); q.pop();
if (up[u] == 0 or l != -d[u]) continue;
   for (int i = 0; i < int(g[u].size()); ++i) {</pre>
    auto e = g[u][i]; int v = e.to;
    auto nd = d[u] + e.c + h[u] - h[v];
    if (e.f <= 0 or d[v] <= nd)</pre>
     continue;
    f[v] = \{u, i\};
    up[v] = min(up[u], e.f);
    q.emplace(-(d[v] = nd), v);
  if (d[T] == INF_C) return nullopt;
  for (size_t i = 0; i < d.size(); i++) h[i]+=d[i];
for (int i = T; i != S; i = f[i].first) {</pre>
   auto &eg = g[f[i].first][f[i].second];
   eg.f -= up[T];
   g[eg.to][eg.r].f += up[T];
  return pair{up[T], h[T]};
public:
MCMF(int n) : g(n), f(n), up(n), d(n, INF_C), h(n) {}
 void add_edge(int s, int t, F c, C w) {
  g[s].emplace_back(t, int(g[t].size()), c, w);
  g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
pair<F, C> solve(int a, int b) {
 F c = 0; C w = 0;
  while (auto r = step(a, b)) {
   c += r->first, w += r->first * r->second;
   fill(d.begin(), d.end(), INF_C);
  return {c, w};
 }
```

4.3 Dinic [ebd802]

```
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
    int to, rev;
    Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
      int u = bfs.front(); bfs.pop();
      for (auto e: G[u]) {
        if (e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
```

```
bfs.push(e.to); lv[e.to] = lv[u] + 1;
     return lv[ed] != -1;
  Cap DFS(int u, Cap f){
     if (u == ed) return f;
     Cap ret = 0;
     for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
       auto &e = G[u][i];
if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
       Cap nf = DFS(e.to, min(f, e.cap));
       ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
       if (f == 0) return ret;
     if (ret == 0) lv[u] = -1;
     return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
   G[u].push_back({v, int(G[v].size()), c});
   G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
  st = st_, ed = ed_; Cap ret = 0;
     while (BFS()) {
       idx.assign(n, 0);
       Cap f = DFS(st, numeric_limits<Cap>::max());
       ret += f;
       if (f == 0) break;
     return ret;
  7
};
```

4.4 Flow Models

- · Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x,y,l,u), connect x o y with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the
 - sum of incoming lower bounds and the sum of outgoing lower
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t\, o\,s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer
 - To minimize, let f be the maximum flow from S to T. Connect $t\,\rightarrow\,s$ with capacity ∞ and let the flow from S to Tbe f'. If $f+f' \neq \sum_{v \in V, in(v)>0}^{\bullet} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
 - 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect $x \to y$ with (cost,cap) = (c,1) if
 - c>0, otherwise connect $y\to x$ with (cost, cap)=(-c,1)3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v) > 0, connect $S \to v$ with (cost, cap) =
 - 5. For each vertex v with d(v) < 0, connect $v \rightarrow T$ with
 - (cost, cap) = (0, -d(v)) 6. Flow from S to T , the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \to v$, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
 - 5. For $v \in G$, connect it with sink $v \to t$ with capacity K + 2T $\left(\sum_{e \in E(v)} w(e)\right) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v~\in~V$ create a copy v' , and connect $u'~\to~v'$ with weight w(u, v).

- 2. Connect v o v' with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
- 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing \widetilde{u} without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.

4.5 General Graph Matching [00732c]

```
namespace matching {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
 for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
 g[u].push_back(v);
 g[v].push_back(u);
int Find(int u) {
 return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
  static int tk = 0; tk++;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
 }
void Blossom(int x, int y, int l) {
 while (Find(x) != l) {
  pre[x] = y, y = match[x];
  if (s[y] == 1) q.push(y), s[y] = 0;
  if (fa[x] == x) fa[x] = l;
  if (fa[y] == y) fa[y] = l;
  x = pre[y];
bool Bfs(int r, int n) {
 for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;
 while (!q.empty()) q.pop();
 q.push(r);
 s[r] = 0;
 while (!q.empty()) {
  int x = q.front(); q.pop();
  for (int u : g[x]) {
   if (s[u] == -1) {
    pre[u] = x, s[u] = 1;
    if (match[u] == n) {
     for (int a = u, b = x, last; b != n; a = last, b =
      last = match[b], match[b] = a, match[a] = b;
     return true;
    q.push(match[u]);
    s[match[u]] = 0;
   } else if (!s[u] && Find(u) != Find(x)) {
    int l = LCA(u, x, n);
    Blossom(x, u, l);
    Blossom(u, x, l);
  }
 return false;
int Solve(int n) {
 for (int x = 0; x < n; ++x) {
  if (match[x] == n) res += Bfs(x, n);
 return res;
```

4.6 Global Min-Cut [1f0306]

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
pair<int, int> phase(int n) {
memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
int s = -1, t = -1;
while (true) {
 int c = -1;
  for (int i = 0; i < n; ++i) {</pre>
   if (del[i] || v[i]) continue;
  if (c == -1 || g[i] > g[c]) c = i;
 if (c == -1) break;
 v[s = t, t = c] = true;
  for (int i = 0; i < n; ++i) {</pre>
  if (del[i] || v[i]) continue;
  g[i] += w[c][i];
return make_pair(s, t);
int mincut(int n) {
int cut = 1e9;
memset(del, false, sizeof(del));
for (int i = 0; i < n - 1; ++i) {</pre>
 int s, t; tie(s, t) = phase(n);
del[t] = true; cut = min(cut, g[t]);
 for (int j = 0; j < n; ++j) {</pre>
  w[s][j] += w[t][j]; w[j][s] += w[j][t];
 }
return cut;
```

4.7 GomoryHu Tree [f8938f]

```
int g[maxn];
vector<edge> GomoryHu(int n){
 vector<edge> rt;
for(int i=1;i<=n;++i)g[i]=1;</pre>
for(int i=2;i<=n;++i){</pre>
  int t=g[i];
 flow.reset(); // clear flows on all edge
  rt.push_back({i,t,flow(i,t)});
 flow.walk(i); // bfs points that connected to i (use
   edges not fully flow)
 for(int j=i+1;j<=n;++j){</pre>
   if(g[j]==t && flow.connect(j))g[j]=i; // check if i
    can reach i
 }
}
return rt;
```

4.8 Kuhn Munkres [1d3c40]

```
class KM {
private:
static constexpr lld INF = 1LL << 60;</pre>
vector<lld> hl,hr,slk;
vector<int> fl,fr,pre,qu;
vector<vector<lld>> w;
vector<bool> vl,vr;
 int n, ql, qr;
bool check(int x) {
 if (vl[x] = true, fl[x] != -1)
   return vr[qu[qr++] = fl[x]] = true;
 while (x != -1) swap(x, fr[fl[x] = pre[x]]);
  return false;
void bfs(int s) {
 fill(slk.begin(), slk.end(), INF);
 fill(vl.begin(), vl.end(), false);
fill(vr.begin(), vr.end(), false);
 ql = qr = 0;
 vr[qu[qr++] = s] = true;
 while (true) {
  lld d;
```

```
while (ql < qr) {</pre>
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
     if(!vl[x]&&slk[x]>=(d=hl[x]+hr[y]-w[x][y])){
      if (pre[x] = y, d) slk[x] = d;
      else if (!check(x)) return;
     }
    }
   }
   d = INF;
   for (int x = 0; x < n; ++x)
    if (!vl[x] && d > slk[x]) d = slk[x];
   for (int x = 0; x < n; ++x) {
    if (vl[x]) hl[x] += d;
    else slk[x] -= d;
    if (vr[x]) hr[x] -= d;
   for (int x = 0; x < n; ++x)
    if (!vl[x] && !slk[x] && !check(x)) return;
 }
public:
void init( int n_ ) {
  qu.resize(n = n_);
 fl.assign(n, -1); fr.assign(n, -1);
 hr.assign(n, 0); hl.resize(n);
 w.assign(n, vector<lld>(n));
 slk.resize(n); pre.resize(n);
  vl.resize(n); vr.resize(n);
 void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 lld solve() {
  for (int i = 0; i < n; ++i)</pre>
  hl[i] = *max_element(w[i].begin(), w[i].end());
  for (int i = 0; i < n; ++i) bfs(i);</pre>
  lld res = 0;
  for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
  return res;
}
} km;
```

4.9 Minimum Cost Circulation [d99194]

```
struct Edge { int to, cap, rev, cost; };
vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
bool mark[kN];
int NegativeCycle(int n) {
memset(mark, false, sizeof(mark));
 memset(dist, 0, sizeof(dist));
 int upd = -1;
 for (int i = 0; i <= n; ++i) {
  for (int j = 0; j < n; ++j) {</pre>
   int idx = 0;
   for (auto &e : g[j]) {
    if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
     dist[e.to] = dist[j] + e.cost;
     pv[e.to] = j, ed[e.to] = idx;
     if (i == n) {
      upd = j;
      while(!mark[upd])mark[upd]=1,upd=pv[upd];
      return upd;
     }
    idx++;
 }
 return -1;
int Solve(int n) {
 int rt = -1, ans = 0;
 while ((rt = NegativeCycle(n)) >= 0) {
  memset(mark, false, sizeof(mark));
  vector<pair<int, int>> cyc;
  while (!mark[rt]) {
   cyc.emplace_back(pv[rt], ed[rt]);
   mark[rt] = true;
  rt = pv[rt];
  reverse(cyc.begin(), cyc.end());
  int cap = kInf;
  for (auto &i : cyc) {
   auto &e = g[i.first][i.second];
```

```
National Taiwan University - ckiseki
   cap = min(cap, e.cap);
                                                                 maxn * 21:
  for (auto &i : cyc) {
   auto &e = g[i.first][i.second];
                                                              queue<int> q;
   e.cap -= cap;
   g[e.to][e.rev].cap += cap;
   ans += e.cost * cap;
                                                                 ] = u; }
return ans:
                                                               slack[x] = 0;
4.10 Minimum Cost Max Flow [6d1b01]
template <typename F, typename C> class MCMF {
 static constexpr F INF_F = numeric_limits<F>::max();
 static constexpr C INF_C = numeric_limits<C>::max();
                                                              void q_push(int x) {
 struct E {
  int to, r;
 F f; C c;
  E() {}
 E(int a, int b, F x, C y)
   : to(a), r(b), f(x), c(y) {}
                                                               st[x] = b;
 };
 vector<vector<E>> g;
 vector<pair<int, int>> f;
 vector<bool> inq;
 vector<F> up; vector<C> d;
 optional<pair<F, C>> step(int S, int T) {
                                                                 [b].begin();
                                                               if (pr % 2 == 1) {
  queue<int> q;
  for (q.push(S), d[S] = 0, up[S] = INF_F;
    not q.empty(); q.pop()) {
   int u = q.front(); inq[u] = false;
   if (up[u] == 0) continue;
                                                               return pr;
   for (int i = 0; i < int(g[u].size()); ++i) {</pre>
    auto e = g[u][i]; int v = e.to;
    if (e.f <= 0 or d[v] <= d[u] + e.c)
                                                               if (u <= n) return;</pre>
     continue:
    d[v] = d[u] + e.c; f[v] = \{u, i\};
                                                               edge e = g[u][v];
    up[v] = min(up[u], e.f);
    if (not inq[v]) q.push(v);
    inq[v] = true;
                                                                 [u][i ^ 1]);
                                                               set_match(xr, v);
   }
  if (d[T] == INF_C) return nullopt;
                                                                 end());
  for (int i = T; i != S; i = f[i].first) {
   auto &eg = g[f[i].first][f[i].second];
                                                               for (; ; ) {
   eg.f -= up[T];
   g[eg.to][eg.r].f += up[T];
                                                                set_match(u, v);
                                                                if (!xnv) return;
  return pair{up[T], d[T]};
public:
MCMF(int n) : g(n), f(n), inq(n), up(n), d(n, INF_C) {}
                                                              }
 void add_edge(int s, int t, F c, C w) {
 g[s].emplace_back(t, int(g[t].size()), c, w);
g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
                                                              static int t = 0;
 pair<F, C> solve(int a, int b) {
 F c = 0; C w = 0;
 while (auto r = step(a, b)) {
                                                                vis[u] = t;
                                                               u = st[match[u]];
   c += r->first, w += r->first * r->second;
   fill(inq.begin(), inq.end(), false);
   fill(d.begin(), d.end(), INF_C);
                                                              return 0;
  return {c, w};
}
                                                              int b = n + 1;
};
      Weighted Matching [60ca53]
                                                               if (b > n_x) ++n_x;
struct WeightGraph {
 static const int inf = INT_MAX;
 static const int maxn = 514;
                                                               flo[b].clear();
 struct edge {
  int u, v, w;
  edge(){}
                                                                 ]]), q_push(y);
  edge(int u, int v, int w): u(u), v(v), w(w) {}
 int n, n_x;
 edge g[maxn \star 2][maxn \star 2];
 int lab[maxn * 2];
                                                                 ]]), q_push(y);
 int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa
```

[maxn * 2];

```
int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
vector<int> flo[maxn * 2];
int e_delta(const edge &e) { return lab[e.u] + lab[e.v
  ] - g[e.u][e.v].w * 2; }
void update_slack(int u, int x) { if (!slack[x] ||
   e_delta(g[u][x]) < e_delta(g[slack[x]][x])) slack[x</pre>
void set_slack(int x) {
for (int u = 1; u <= n; ++u)</pre>
 if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
  update_slack(u, x);
if (x \le n) q.push(x);
else for (size_t i = 0; i < flo[x].size(); i++)</pre>
   q_push(flo[x][i]);
void set_st(int x, int b) {
if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)</pre>
    set_st(flo[x][i], b);
int get_pr(int b, int xr) {
int pr = find(flo[b].begin(), flo[b].end(), xr) - flo
 reverse(flo[b].begin() + 1, flo[b].end());
 return (int)flo[b].size() - pr;
void set_match(int u, int v) {
match[u] = g[u][v].v;
 int xr = flo_from[u][e.u], pr = get_pr(u, xr);
for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo</pre>
rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
void augment(int u, int v) {
 int xnv = st[match[u]];
 set_match(xnv, st[pa[xnv]]);
 u = st[pa[xnv]], v = xnv;
int get_lca(int u, int v) {
for (++t; u || v; swap(u, v)) {
 if (u == 0) continue;
 if (vis[u] == t) return u;
 if (u) u = st[pa[u]];
void add_blossom(int u, int lca, int v) {
while (b <= n_x && st[b]) ++b;</pre>
lab[b] = 0, S[b] = 0;
match[b] = match[lca];
flo[b].push_back(lca);
for (int x = u, y; x != lca; x = st[pa[y]])
 flo[b].push_back(x), flo[b].push_back(y = st[match[x
 reverse(flo[b].begin() + 1, flo[b].end());
for (int x = v, y; x != lca; x = st[pa[y]])
 flo[b].push_back(x), flo[b].push_back(y = st[match[x
set_st(b, b);
```

```
for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].w
 for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;</pre>
 for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  for (int x = 1; x <= n_x; ++x)</pre>
   if (g[b][x].w == 0 || e_delta(g[xs][x]) < e_delta(g</pre>
   [b][x]))
  g[b][x] = g[xs][x], g[x][b] = g[x][xs];

for (int x = 1; x <= n; ++x)
   if (flo_from[xs][x]) flo_from[b][x] = xs;
 set_slack(b);
}
void expand_blossom(int b) {
 for (size_t i = 0; i < flo[b].size(); ++i)
    set_st(flo[b][i], flo[b][i]);</pre>
 int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
   xr);
 for (int i = 0; i < pr; i += 2) {</pre>
  int xs = flo[b][i], xns = flo[b][i + 1];
  pa[xs] = g[xns][xs].u;
  S[xs] = 1, S[xns] = 0;
slack[xs] = 0, set_slack(xns);
  q_push(xns);
 S[xr] = 1, pa[xr] = pa[b];
 for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  S[xs] = -1, set_slack(xs);
 st[b] = 0;
bool on_found_edge(const edge &e) {
 int u = st[e.u], v = st[e.v];
 if (S[v] == -1) {
  pa[v] = e.u, S[v] = 1;
  int nu = st[match[v]];
  slack[v] = slack[nu] = 0;
  S[nu] = 0, q_push(nu);
 } else if (S[v] == 0) {
  int lca = get_lca(u, v);
if (!lca) return augment(u,v), augment(v,u), true;
  else add_blossom(u, lca, v);
 }
 return false;
bool matching() {
 memset(S + \overline{1}, -1, sizeof(int) * n_x);
 memset(slack + 1, 0, sizeof(int) * n_x);
 q = queue<int>();
 for (int x = 1; x <= n_x; ++x)
  if (st[x] == x && :match[x]) pa[x] = 0, S[x] = 0,
   q_push(x);
 if (q.empty()) return false;
 for (; ; ) {
  while (q.size()) {
   int u = q.front(); q.pop();
   if (S[st[u]] == 1) continue;
   for (int v = 1; v <= n; ++v)</pre>
    if (g[u][v].w > 0 && st[u] != st[v]) {
     if (e_delta(g[u][v]) == 0) {
      if (on_found_edge(g[u][v])) return true;
     } else update_slack(u, st[v]);
    }
  int d = inf;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b \&\& S[b] == 1) d = min(d, lab[b] / 2)
  for (int x = 1; x \le n_x; ++x)
   if (st[x] == x && slack[x]) {
    if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x])
   );
    else if (S[x] == 0) d = min(d, e_delta(g[slack[x
   ]][x]) / 2);
  for (int u = 1; u <= n; ++u) {</pre>
   if (S[st[u]] == 0) {
    if (lab[u] <= d) return 0;</pre>
    lab[u] -= d;
   } else if (S[st[u]] == 1) lab[u] += d;
```

```
for (int b = n + 1; b \le n_x; ++b)
    if (st[b] == b) {
     if (S[st[b]] == 0) lab[b] += d * 2;
     else if (S[st[b]] == 1) lab[b] -= d * 2;
   q = queue<int>();
   for (int x = 1; x \le n_x; ++x)
    if (st[x] == x && slack[x] && st[slack[x]] != x &&
     e_delta(g[slack[x]][x]) == 0)
     if (on_found_edge(g[slack[x]][x])) return true;
   for (int b = n + 1; b <= n_x; ++b)
if (st[b] == b && S[b] == 1 && lab[b] == 0)</pre>
     expand_blossom(b);
  return false;
 pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n_x = n;
  int n_matches = 0;
  long long tot_weight = 0;
  for (int u = 0; u <= n; ++u) st[u] = u, flo[u].clear
     ();
  int w_max = 0;
  for (int u = 1; u <= n; ++u)</pre>
   for (int v = 1; v \le n; ++v) {
    flo_from[u][v] = (u == v ? u : 0);
    w_max = max(w_max, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
  while (matching()) ++n_matches;
  for (int u = 1; u <= n; ++u)</pre>
   if (match[u] && match[u] < u)</pre>
    tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
 void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
  g[vi][ui].w = wi; }
 void init(int _n) {
  n = _n;
  for (int u = 1; u <= n; ++u)</pre>
   for (int v = 1; v <= n; ++v)
    g[u][v] = edge(u, v, 0);
 }
};
```

5 Math

5.1 Common Bounds

$$\begin{split} p(0) &= 1, \; p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n-k(3k-1)/2) \\ & p(n) \approx 0.145/n \cdot \exp(2.56\sqrt{n}) \\ \\ & \frac{n}{\max_{i \leq n} (d(i))} \mid 12 \quad 32 \quad 240 \quad 1344 \quad 6720 \quad 26880 \quad 103680 \\ \\ & \frac{n}{\binom{2n}{n}} \mid 2 \quad 6 \quad 26 \quad 20 \quad 70 \quad 252 \quad 924 \quad 3432 \quad 12870 \quad 48620 \quad 184756 \end{split}$$

5.2 Strling Number

First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1)\dots(x+n-1) = \sum_{k=0}^n S_1(n,k)x^k$$

$$g(x) = x(x+1)\dots(x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into knonempty sets.

```
S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)
```

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

ax+by=gcd [d0cbdd]

```
// ax+ny = 1, ax+ny == ax == 1 \ (mod \ n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
if (y == 0) g=x,a=1,b=0;
else exgcd(y,x\%y,g,b,a),b=(x/y)*a;
```

5.4 Chinese Remainder [d69e74]

```
// please ensure r_i\in[0,m_i)
bool crt(lld &m1, lld &r1, lld m2, lld r2) {
  if (m2 > m1) swap(m1, m2), swap(r1, r2);
  lld g, a, b; exgcd(m1, m2, g, a, b);
  if ((r2 - r1) % g != 0) return false;
m2 /= g; lld D = (r2 - r1) / g % m2 * a % m2;
  r1 += (D < 0 ? D + m2 : D) * m1; m1 *= m2;
  assert (r1 >= 0 \&\& r1 < m1);
  return true;
```

5.5 De-Bruijn [7f536e]

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
 if (t > n) {
 if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
} else {
  aux[t] = aux[t - p];
 db(t + 1, p, n, k);

for (int i = aux[t - p] + 1; i < k; ++i) {
   aux[t] = i;
  db(t + 1, t, n, k);
 }
int de_bruijn(int k, int n) {
// return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
if (k == 1) {
 res[0] = 0;
 return 1;
for (int i = 0; i < k * n; i++) aux[i] = 0;
sz = 0;
db(1, 1, n, k);
return sz;
```

5.6 DiscreteLog [86e463]

```
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
// x^? \setminus equiv y \pmod{M}
Int t = 1, c = 0, g = 1;

for (Int M_ = M; M_ > 0; M_ >>= 1)
 g = g * x % M;
for (g = gcd(g, M); t % g != 0; ++c) {
 if (t == y) return c;
 t = t * x % M;
if (y % g != 0) return -1;
t /= g, y /= g, M /= g;
Int h = 0, gs = 1;
for (; h * h < M; ++h) gs = gs * x % M;
unordered_map<Int, Int> bs;
 for (Int s = 0; s < h; bs[y] = ++s)
 y = y * x % M;
for (Int s = 0; s < M; s += h) {</pre>
 t = t * gs % M;
  if (bs.count(t)) return c + s + h - bs[t];
return -1;
```

5.7 Quadratic residue [leabad]

```
int get_root(int n, int P) { // ensure 0 <= n < p</pre>
 if (P == 2 or n == 0) return n;
 auto check = [&](int x) {
 return modpow(x, (P - 1) / 2, P); };
if (check(n) != 1) return -1;
 mt19937 \text{ rnd}(7122); \text{ lld } z = 1, w;
 while (check(w = (z * z - n + P) % P) != P - 1)
  z = rnd() \% P;
 const auto M = [P, w](auto &u, auto &v) {
  auto [a, b] = u; auto [c, d] = v;
return make_pair((a * c + b * d % P * w) % P,
     (a * d + b * c) % P);
 };
 pair<lld, lld> r(1, 0), e(z, 1);
for (int w = (P + 1) / 2; w; w >>= 1, e = M(e, e))
 if (w & 1) r = M(r, e);
 return r.first; // sqrt(n) mod P where P is prime
```

5.8 Extended Euler

```
a^b \equiv \begin{cases} a^{(b \mod \varphi(m)) + \varphi(m)} & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases}
```

5.9 Extended FloorSum

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \end{cases} \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1)) - h(c, c-b-1, a, m-1), & \text{otherwise} \end{cases} \\ h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{c} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

5.10 FloorSum [bda6b2]

```
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
// @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
 llu ans = 0;
 while (true)
  if (a >= m) {
   ans += n * (n - 1) / 2 * (a / m); a %= m;
  if (b >= m) {
   ans += n * (b / m); b \% = m;
  llu y_max = a * n + b;
  if (y_max < m) break;</pre>
  // y_max < m * (n + 1)
// floor(y_max / m) <= n
  n = (llu)(y_max / m), b = (llu)(y_max % m);
  swap(m, a);
 return ans;
lld floor_sum(lld n, lld m, lld a, lld b) {
 if (a < 0) {
  llu a2 = (a \% m + m) \% m;
 ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
  a = a2;
 if (b < 0) {
  llu b2 = (b \% m + m) \% m;
  ans -= 1ULL * n * ((b2 - b) / m);
  b = b2;
 return ans + floor_sum_unsigned(n, m, a, b);
```

5.11 ModMin [07d5e1]

5.12 Fast Fourier Transform [993ee3]

```
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
for (int i = 0; i <= maxn; i++)</pre>
 omega[i] = cplx(cos(2 * pi * j / maxn),
     sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
int z = __builtin_ctz(n) - 1;
for (int i = 0; i < n; ++i) {</pre>
 int x = 0, j = 0;
 for (;(1 << j) < n;++j) x^=(i >> j & 1) << (z - j);
 if (x > i) swap(v[x], v[i]);
for (int s = 2; s <= n; s <<= 1) {
 int z = s >> 1;
  for (int i = 0; i < n; i += s) {</pre>
  for (int k = 0; k < z; ++k) {
  cplx x = v[i + z + k] * omega[maxn / s * k];</pre>
   v[i + z + k] = v[i + k] - x;
    v[i + k] = v[i + k] + x;
   }
void ifft(vector<cplx> &v, int n) {
fft(v, n); reverse(v.begin() + 1, v.end());
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
VL convolution(const VI &a, const VI &b) {
// Should be able to handle N <= 10^5, C <= 10^4
int sz = 1;
while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
vector<cplx> v(sz);
for (int i = 0; i < sz; ++i) {</pre>
 double re = i < a.size() ? a[i] : 0;
 double im = i < b.size() ? b[i] : 0;</pre>
 v[i] = cplx(re, im);
fft(v, sz);
for (int i = 0; i <= sz / 2; ++i) {</pre>
 int j = (sz - i) & (sz - 1);
 cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
    * cplx(0, -0.25);
 if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
    ].conj()) * cplx(0, -0.25);
 v[i] = x;
ifft(v, sz);
VL c(sz);
for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
VI convolution_mod(const VI &a, const VI &b, int p) {
int sz = 1;
while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
vector<cplx> fa(sz), fb(sz);
for (int i = 0; i < (int)a.size(); ++i)</pre>
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
for (int i = 0; i < (int)b.size(); ++i)</pre>
 fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
fft(fa, sz), fft(fb, sz);
double r = 0.25 / sz;
cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
```

```
for (int i = 0; i <= (sz >> 1); ++i) {
  int j = (sz - i) & (sz - 1);
  cplx a1 = (fa[i] + fa[j].conj());
  cplx a2 = (fa[i] - fa[j].conj()) * r2;
  cplx b1 = (fb[i] + fb[j].conj()) * r3;
  cplx b2 = (fb[i] - fb[j].conj()) * r4;
  if (i != j) {
   cplx c1 = (fa[j] + fa[i].conj());
   cplx c2 = (fa[j] - fa[i].conj()) * r2;
   cplx d1 = (fb[j] + fb[i].conj()) * r3;
   cplx d2 = (fb[j] - fb[i].conj()) * r4;
   fa[i] = c1 * d1 + c2 * d2 * r5;
   fb[i] = c1 * d2 + c2 * d1;
  fa[j] = a1 * b1 + a2 * b2 * r5;
  fb[j] = a1 * b2 + a2 * b1;
 fft(fa, sz), fft(fb, sz);
 vector<int> res(sz);
 for (int i = 0; i < sz; ++i) {</pre>
  long long a = round(fa[i].re), b = round(fb[i].re),
       c = round(fa[i].im);
  res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
 return res;
}}
5.13 FWT [c5167a]
/* or convolution:
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
 * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
void fwt(int x[], int N, bool inv = false) {
 for (int d = 1; d < N; d <<= 1)
  for (int s = 0; s < N; s += d * 2)
   for (int i = s; i < s + d; i++) {
    int j = i + d, ta = x[i], tb = x[j];
    x[i] = modadd(ta, tb);
    x[j] = modsub(ta, tb);
 if (inv) {
  const int invn = modinv(N);
  for (int i = 0; i < N; i++)
   x[i] = modmul(x[i], invn);
5.14 CRT for arbitrary mod [7272c4]
const int mod = 1000000007;
const int M1 = 985661441; // G = 3 for M1, M2, M3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
  static_assert (M1 <= M2 && M2 <= M3);</pre>
  constexpr int64_t r12 = modpow(M1, M2-2, M2);
  constexpr int64_t r13 = modpow(M1, M3-2, M3);
  constexpr int64_t r23 = modpow(M2, M3-2, M3);
  constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
  B = (B - A + M2) * r12 % M2;
  C = (C - A + M3) * r13 % M3;
  C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
5.15 NTT [946e8e]
```

```
template <int mod, int G, int maxn> struct NTT {
    static_assert (maxn == (maxn & -maxn));
    int roots[maxn];
NTT () {
    int r = modpow(G, (mod - 1) / maxn);
    for (int i = maxn >> 1; i; i >>= 1) {
        roots[i] = 1;
        for (int j = 1; j < i; j++)
            roots[i + j] = modmul(roots[i + j - 1], r);
        r = modmul(r, r);
    }
}
// n must be 2^k, and 0 <= F[i] < mod
void operator()(int F[], int n, bool inv = false) {
    for (int i = 0, j = 0; i < n; i++) {
        if (i < j) swap(F[i], F[j]);
    }
}</pre>
```

```
for (int k = n > 1; (j^k < k; k > = 1);
  for (int s = 1; s < n; s *= 2) {
   for (int i = 0; i < n; i += s * 2) {</pre>
    for (int j = 0; j < s; j++) {
     int a = F[i+j];
     int b = modmul(F[i+j+s], roots[s+j]);
     F[i+j] = modadd(a, b); // a + b
     F[i+j+s] = modsub(a, b); // a - b
  }
 if (inv) {
   int invn = modinv(n);
  for (int i = 0; i < n; i++)</pre>
   F[i] = modmul(F[i], invn);
   reverse(F + 1, F + n);
 }
}
};
```

5.16 Partition Number [9bb845]

```
ans[0] = tmp[0] = 1;
for (int i = 1; i * i <= n; i++) {
  for (int rep = 0; rep < 2; rep++)
    for (int j = i; j <= n - i * i; j++)
    modadd(tmp[j], tmp[j-i]);
  for (int j = i * i; j <= n; j++)
    modadd(ans[j], tmp[j - i * i]);
}</pre>
```

5.17 Pi Count (+Linear Sieve) [47e0de]

```
static constexpr int N = 1000000 + 5;
lld pi[N];
vector<int> primes;
bool sieved[N];
lld cube_root(lld x){
 lld s=cbrt(x-static_cast<long double>(0.1));
 while(s*s*s <= x) ++s;
return s-1;
lld square_root(lld x){
 lld s=sqrt(x-static_cast<long double>(0.1));
 while(s*s <= x) ++s;
 return s-1;
void init(){
 primes.reserve(N);
 primes.push_back(1);
 for(int i=2;i<N;i++) {</pre>
  if(!sieved[i]) primes.push_back(i);
 pi[i] = !sieved[i] + pi[i-1];
 for(int p: primes) if(p > 1) {
   if(p * i >= N) break;
   sieved[p * i] = true;
   if(p % i == 0) break;
}
lld phi(lld m, lld n) {
 static constexpr int MM = 80000, NN = 500;
 static lld val[MM][NN];
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
 lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
 if(m<MM&&n<NN) val[m][n] = ret+1;
 return ret:
lld pi_count(lld);
lld P2(lld m, lld n) {
lld sm = square_root(m), ret = 0;
 for(lld i = n+1;primes[i]<=sm;i++)</pre>
 ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
 return ret;
lld pi_count(lld m) {
 if(m < N) return pi[m];</pre>
 lld n = pi_count(cube_root(m));
 return phi(m, n) + n - 1 - P2(m, n);
```

5.18 Miller Rabin [0edab2]

```
bool isprime(llu x) {
 static auto witn = [](llu a, llu n, int t) {
  if (!a) return false;
  while (t--) {
   llu a2 = mmul(a, a, n);
   if (a2 == 1 && a != 1 && a != n - 1) return true;
  return a != 1:
 if (x < 2) return false;</pre>
 if (!(x & 1)) return x == 2;
 int t = __builtin_ctzll(x - 1);
 llu odd = (x - 1) >> t;
 for (llu m:
  {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
  if (witn(mpow(m % x, odd, x), x, t))
   return false;
 return true;
```

5.19 Pollard Rho [2aclad]

```
// does not work when n is prime
// return any non-trivial factor
llu pollard_rho(llu n) {
 static auto f = [](llu x, llu k, llu m) {
    return add(k, mul(x, x, m), m); };
 if (!(n & 1)) return 2;
 mt19937 rnd(120821011);
 while (true) {
  llu y = 2, yy = y, x = rnd() % n, t = 1;
  for (llu sz = 2; t == 1; sz <<= 1, y = yy) {
   for (llu i = 0; t == 1 && i < sz; ++i) {
    yy = f(yy, x, n);
    t = gcd(yy > y ? yy - y : y - yy, n);
  if (t != 1 && t != n) return t;
 }
}
```

5.20 Berlekamp Massey [a94d00]

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
 vector<T> d(output.size() + 1), me, he;
for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
  for (size_t j = 0; j < me.size(); ++j)</pre>
   d[i] += output[i - j - 2] * me[j];
  if ((d[i] -= output[i - 1]) == 0) continue;
  if (me.empty()) {
   me.resize(f = i);
   continue;
  vector<T> o(i - f - 1);
  T k = -d[i] / d[f]; o.push_back(-k);
  for (T x : he) o.push_back(x * k);
  if (o.size() < me.size()) o.resize(me.size());</pre>
  for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
  if (i-f+he.size() >= me.size()) he = me, f = i;
  me = o;
 }
 return me;
```

5.21 Charateristic Polynomial [e006eb]

```
#define rep(x, y, z) for (int x=y; x<z; x++)
using VI = vector<int>; using VVI = vector<VI>;
void Hessenberg(VVI &H, int N) {
    for (int i = 0; i < N - 2; ++i) {
        for (int j = i + 1; j < N; ++j) if (H[j][i]) {
            rep(k, i, N) swap(H[i+1][k], H[j][k]);
            rep(k, 0, N) swap(H[k][i+1], H[k][j]);
            break;
        }
        if (!H[i + 1][i]) continue;
        for (int j = i + 2; j < N; ++j) {
            int co = mul(modinv(H[i + 1][i]), H[j][i]);
            rep(k, i, N) subeq(H[j][k], mul(H[i+1][k], co));
            rep(k, 0, N) addeq(H[k][i+1], mul(H[k][j], co));
        }
}</pre>
```

```
}
VI CharacteristicPoly(VVI &A) {
int N = A.size(); Hessenberg(A, N);
VVI P(N + 1, VI(N + 1)); P[0][0] = 1;
for (int i = 1; i <= N; ++i) {</pre>
 rep(j, 0, i+1) P[i][j] = j ? P[i-1][j-1] : 0;
  for (int j = i - 1, val = 1; j >= 0; --j) {
   int co = mul(val, A[j][i - 1]);
   rep(k, 0, j+1) subeq(P[i][k], mul(P[j][k], co));
   if (j) val = mul(val, A[j][j - 1]);
 }
if (N & 1) for (int &pi: P[N]) pi = sub(0, pi);
return P[N]; // test: 2021 PTZ Korea K
```

5.22 Polynomial Operations [d40491]

```
using V = vector<int>;
#define fi(l, r) for (int i = int(l); i < int(r); ++i)
template <int mod, int G, int maxn> struct Poly : V {
static uint32_t n2k(uint32_t n) {
 if (n <= 1) return 1;
 return 1u << (32 - __builtin_clz(n - 1));</pre>
}
static NTT<mod,G,maxn> ntt; // coefficients in [0, P)
explicit Poly(int n = 1) : V(n) {}
Poly(const V &v) : V(v) {}
Poly(const Poly &p, size_t n) : V(n) {
 copy_n(p.data(), min(p.size(), n), data());
Poly &irev() { return reverse(data(), data() + size())
    , *this; }
Poly &isz(int sz) { return resize(sz), *this; }
Poly &iadd(const Poly &rhs) { // n() == rhs.n()
 fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
 return *this;
Poly &imul(int k) {
 fi(0, size())(*this)[i] = modmul((*this)[i], k);
 return *this;
Poly Mul(const Poly &rhs) const {
 const int sz = n2k(size() + rhs.size() - 1);
 Poly X(*this, sz), Y(rhs, sz);
 ntt(X.data(), sz), ntt(Y.data(), sz);
 fi(0, sz) X[i] = modmul(X[i], Y[i]);
 ntt(X.data(), sz, true);
 return X.isz(size() + rhs.size() - 1);
Poly Inv() const { // coef[0] != 0
 if (size() == 1) return V{modinv(*begin())};
 const int sz = n2k(size() * 2);
 Poly X = Poly(*this, (size() + 1) / 2).Inv().isz(sz),
     Y(*this, sz);
  ntt(X.data(), sz), ntt(Y.data(), sz);
  fi(0, sz) X[i] = modmul(X[i], modsub(2, modmul(X[i],
    Y[i])));
 ntt(X.data(), sz, true);
 return X.isz(size());
Poly Sqrt() const { // coef[0] \in [1, mod)^2
 if (size() == 1) return V{QuadraticResidue((*this))
    [0], mod)};
  Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
    size());
 return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
    + 1);
pair<Poly, Poly> DivMod(const Poly &rhs) const {
 if (size() < rhs.size()) return {V{0}, *this};</pre>
 const int sz = size() - rhs.size() + 1;
 Poly X(rhs); X.irev().isz(sz);
 Poly Y(*this); Y.irev().isz(sz);
 Poly Q = Y.Mul(X.Inv()).isz(sz).irev();
 X = rhs.Mul(Q), Y = *this;
fi(0, size()) Y[i] = modsub(Y[i], X[i]);
 return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
Poly Dx() const {
 Poly ret(size() - 1);
```

```
fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
        17):
    return ret.isz(max<int>(1, ret.size()));
  Poly Sx() const {
    Poly ret(size() + 1);
    fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
        this)[i]);
  Poly Ln() const { // coef[0] == 1; res[0] == 0
    return Dx().Mul(Inv()).Sx().isz(size());
  Poly Exp() const { // coef[0] == 0; res[0] == 1
    if (size() == 1) return V{1};
    Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size)
         ());
    Poly Y = X.Ln(); Y[0] = mod - 1;
    fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
    return X.Mul(Y).isz(size());
  Poly Pow(const string &K) const {
    int nz = 0;
    while (nz < size() && !(*this)[nz]) ++nz;</pre>
    int nk = 0, nk2 = 0;
    for (char c : K) {
     nk = (nk * 10 + c - '0') \% mod;
      nk2 = nk2 * 10 + c - '0';
      if (nk2 * nz >= size())
        return Poly(size());
      nk2 %= mod - 1;
    if (!nk && !nk2) return Poly(V{1}, size());
    Poly X = V(data() + nz, data() + size() - nz * (nk2 - nz)
          1));
    int x0 = X[0];
    return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
        modpow(x0, nk2)).irev().isz(size()).irev();
  V Eval(V x) const {
    if (x.empty()) return {};
    const size_t n = max(x.size(), size());
    vector<Poly> t(n * 2, V{1, 0}), f(n * 2);
    for (size_t i = 0; i < x.size(); ++i)</pre>
     t[n + i] = V{1, mod-x[i]};
    for (size_t i = n - 1; i > 0; --i)
      t[i] = t[i * 2].Mul(t[i * 2 + 1]);
    f[1] = Poly(*this, n).irev().Mul(t[1].Inv()).isz(n).
        irev();
    for (size_t i = 1; i < n; ++i) {</pre>
      auto o = f[i]; auto sz = o.size();
      f[i*2] = o.irev().Mul(t[i*2+1]).isz(sz).irev().isz(t
         [i*2].size());
      f[i*2+1] = o.Mul(t[i*2]).isz(sz).irev().isz(t[i*2]).isz(sz).irev().isz(t[i*2]).isz(sz).irev().isz(t[i*2]).isz(sz).irev().isz(t[i*2]).isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().isz(sz).irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev().irev(
         *2+1].size());
    for (size_t i=0;i<x.size();++i) x[i] = f[n+i][0];</pre>
    return x;
  static int LinearRecursion(const V &a, const V &c,
        int64_t n) { // a_n = \sum_{j=0}^{n} a_{j}(n-j)
    const int k = (int)a.size();
    assert((int)c.size() == k + 1);
    Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
    fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
    C[k] = 1;
    while (n) {
      if (n % 2) W = W.Mul(M).DivMod(C).second;
     n /= 2, M = M.Mul(M).DivMod(C).second;
    int ret = 0;
    fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
    return ret;
  }
#undef fi
using Poly_t = Poly<998244353, 3, 1 << 20>;
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
5.23 Simplex [e975d5]
```

namespace simplex {

```
// maximize c^Tx under Ax <= B
                                                                 #define IM imag
// return VD(n, -inf) if the solution doesn't exist
                                                                 #define RE real
// return VD(n, +inf) if the solution is unbounded
                                                                 using lld = int64_t;
                                                                 using llf = long double;
using VD = vector<double>;
using VVD = vector<vector<double>>;
                                                                using PT = std::complex<lld>;
                                                                using PTF = std::complex<llf>;
const double eps = 1e-9;
const double inf = 1e+9;
                                                                 using P = PT;
int n, m;
                                                                 llf abs(P p) { return sqrtl(norm(p)); }
                                                                PTF toPTF(PT p) { return PTF{RE(p), IM(p)}; }
int sgn(lld x) { return (x > 0) - (x < 0); }
lld dot(P a, P b) { return RE(conj(a) * b); }</pre>
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
 double inv = 1.0 / d[r][s];
                                                                 lld cross(P a, P b) { return IM(conj(a) * b); }
                                                                 int ori(P a, P b, P c) {
 for (int i = 0; i < m + 2; ++i)</pre>
  for (int j = 0; j < n + 2; ++j)
                                                                  return sgn(cross(b - a, c - a));
   if (i != r && j != s)
                                                                 int quad(P p) {
    d[i][j] -= d[r][j] * d[i][s] * inv;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;</pre>
                                                                  return (IM(p) == 0) // use sgn for PTF
 for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
                                                                   ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
 d[r][s] = inv; swap(p[r], q[s]);
                                                                 int argCmp(P a, P b) {
                                                                  // returns 0/+-1, starts from theta = -PI
bool phase(int z) {
 int x = m + z;
                                                                  int qa = quad(a), qb = quad(b);
 while (true) {
                                                                  if (qa != qb) return sgn(qa - qb);
  int s = -1;
                                                                  return sgn(cross(b, a));
  for (int i = 0; i <= n; ++i) {</pre>
   if (!z && q[i] == -1) continue;
                                                                 template <typename V> llf area(const V & pt) {
   if (s == -1 || d[x][i] < d[x][s]) s = i;
                                                                  lld ret = 0;
                                                                  for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
  if (d[x][s] > -eps) return true;
                                                                  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
  int r = -1;
                                                                  return ret / 2.0;
  for (int i = 0; i < m; ++i) {</pre>
   if (d[i][s] < eps) continue;</pre>
                                                                P rot90(P p) { return P{-IM(p), RE(p)}; }
                                                                PTF project(PTF p, PTF q) { // p onto q
   if (r == -1 ||
    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
                                                                 return dot(p, q) * q / dot(q, q); // dot<llf>
  if (r == -1) return false;
                                                                 6.2 2D Convex Hull [ecba37]
  pivot(r, s);
 }
                                                                   from NaCl, counterclockwise, be careful of n<=2
                                                                 vector<P> convex_hull(vector<P> v) {
VD solve(const VVD &a, const VD &b, const VD &c) {
                                                                  sort(all(v)); // by X then Y
 m = b.size(), n = c.size();
                                                                  if (v[0] == v.back()) return {v[0]};
 d = VVD(m + 2, VD(n + 2));
                                                                  int t = 0, s = 1; vector<P> h(v.size() + 1);
 for (int i = 0; i < m; ++i)</pre>
                                                                  for (int _ = 2; _--; s = t--, reverse(all(v)))
 for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
                                                                   for (P p : v) {
 p.resize(m), q.resize(n + 1);
                                                                    while (t>s && ori(p, h[t-1], h[t-2]) >= 0) t--;
 for (int i = 0; i < m; ++i)</pre>
                                                                    h[t++] = p;
 p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
 for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];</pre>
                                                                  return h.resize(t), h;
 q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
 for (int i = 1; i < m; ++i)</pre>
                                                                 6.3 2D Farthest Pair [ceb2ae]
  if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
                                                                 // p is CCW convex hull w/o colinear points
 if (d[r][n + 1] < -eps) {</pre>
                                                                 int n = p.size(), pos = 1; lld ans = 0;
  pivot(r, n);
                                                                 for (int i = 0; i < n; i++) {</pre>
  if (!phase(1) || d[m + 1][n + 1] < -eps)</pre>
                                                                  P = p[(i + 1) \% n] - p[i];
   return VD(n, -inf);
                                                                  while (cross(e, p[(pos + 1) % n] - p[i]) >
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
                                                                      cross(e, p[pos] - p[i]))
   int s = min_element(d[i].begin(), d[i].end() - 1)
                                                                   pos = (pos + 1) % n;
        - d[i].begin();
                                                                  for (int j: {i, (i + 1) % n})
   pivot(i, s);
                                                                   ans = max(ans, norm(p[pos] - p[j]));
  }
                                                                } // tested @ AOJ CGL_4_B
 if (!phase(0)) return VD(n, inf);
                                                                       MinMax Enclosing Rect [c66dbf]
 VD x(n);
 for (int i = 0; i < m; ++i)</pre>
                                                                 // from 8BQube, plz ensure p is strict convex hull
  if (p[i] < n) x[p[i]] = d[i][n + 1];</pre>
                                                                 const llf INF = 1e18, qi = acos(-1) / 2 * 3;
 return x;
                                                                 pair<llf, llf> solve(vector<P> &p) {
                                                                 #define Z(v) (p[v] - p[i])
                                                                  llf mx = 0, mn = INF;
5.24 Simplex Construction
                                                                  int n = (int)p.size(); p.emplace_back(p[0]);
Standard form: maximize \sum_{1 \le i \le n} c_i x_i such that for all 1 \le j \le m,
                                                                  for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {
\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j and x_i \geq 0 for all 1 \leq i \leq n.
                                                                   P e = Z(i + 1);
  1. In case of minimization, let c_i'=-c_i
2. \sum_{1\leq i\leq n}A_{ji}x_i\geq b_j\to\sum_{1\leq i\leq n}-A_{ji}x_i\leq -b_j
                                                                   while (cross(e, Z(u + 1)) > cross(e, Z(u)))
                                                                    u = (u + 1) \% n;
  3. \sum_{1 \leq i \leq n} A_{ji} x_i = b_j
                                                                   while (dot(e, Z(r + 1)) > dot(e, Z(r)))
                                                                    r = (r + 1) \% n;
       • \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
                                                                   if (!i) l = (r + 1) % n;
        • \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j
                                                                   while (dot(e, Z(l + 1)) < dot(e, Z(l)))
  4. If x_i has no lower bound, replace x_i with x_i - x_i^\prime
                                                                    l = (l + 1) \% n;
                                                                   P D = p[r] - p[l];
                                                                   mn = min(mn, dot(e, D) / llf(norm(e)) * cross(e, Z(u)
```

));

Geometry

6.1 Basic Geometry [17fa9b]

```
llf B = sqrt(norm(D)) * sqrt(norm(Z(u)));
llf deg = (qi - acos(dot(D, Z(u)) / B)) / 2;
mx = max(mx, B * sin(deg) * sin(deg));
}
return {mn, mx};
}
```

6.5 Minkowski Sum [c71bec]

```
// A, B are convex hull rotate to min by (X, Y)
vector<P> Minkowski(vector<P> A, vector<P> B) {
  vector<P> C(1, A[0] + B[0]), s1, s2;
  const int N = (int)A.size(), M = (int)B.size();
  for(int i = 0; i < N; ++i)
    s1.pb(A[(i + 1) % N] - A[i]);
  for(int i = 0; i < M; i++)
    s2.pb(B[(i + 1) % M] - B[i]);
  for(int i = 0, j = 0; i < N || j < M;)
    if (j >= N || (i < M && cross(s1[i], s2[j]) >= 0))
        C.pb(C.back() + s1[i++]);
  else
        C.pb(C.back() + s2[j++]);
  return hull(C), C;
}
```

6.6 Segment Intersection [60d016]

```
struct Seg { // closed segment
P st, dir; // represent st + t*dir for 0 <= t <= 1
Seg(P s, P e) : st(s), dir(e - s) {}
 static bool valid(lld p, lld q) {
  // is there t s.t. 0 <= t <= 1 && qt == p ?
  if (q < 0) q = -q, p = -p;
  return 0 <= p && p <= q;
 vector<P> ends() const { return { st, st + dir }; }
template <typename T> bool isInter(T A, P p) {
 if (A.dir == P(0)) return p == A.st; // BE CAREFUL
 return cross(p - A.st, A.dir) == 0 &&
  T::valid(dot(p - A.st, A.dir), norm(A.dir));
template <typename U, typename V>
bool isInter(U A, V B) {
   if (cross(A.dir, B.dir) == 0) { // BE CAREFUL
 bool res = false;
  for (P p: A.ends()) res |= isInter(B, p);
  for (P p: B.ends()) res |= isInter(A, p);
 return res:
 P D = B.st - A.st; lld C = cross(A.dir, B.dir);
 return U::valid(cross(D, B.dir), C) &&
  V::valid(cross(D, A.dir), C);
```

6.7 Half Plane Intersection [e98068]

```
struct Line {
 P st, ed, dir;
Line (P s, P e) : st(s), ed(e), dir(e - s) {}
}; using L = const Line &;
PTF intersect(L A, L B) {
 llf t = cross(B.st - A.st, B.dir) /
  llf(cross(A.dir, B.dir));
 return toPTF(A.st) + toPTF(A.dir) * t; // C^3 / C^2
bool cov(L l, L A, L B) {
i128 u = cross(B.st-A.st, B.dir);
i128 v = cross(A.dir, B.dir);
 // ori(l.st, l.ed, A.st + A.dir*(u/v)) <= 0?
 i128 x = RE(A.dir) * u + RE(A.st - l.st) * v;
 i128 y = IM(A.dir) * u + IM(A.st - l.st) * v;
return sgn(x*IM(l.dir) - y*RE(l.dir)) * sgn(v) >= 0;
} // x, y are C^3, also sgn<i128> is needed
bool operator<(L a, L b) {</pre>
 if (int c = argCmp(a.dir, b.dir)) return c == -1;
 return ori(a.st, a.ed, b.st) < 0;</pre>
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
// the half plane is the LHS when going from st to ed
llf HPI(vector<Line> &q) {
 sort(q.begin(), q.end());
 int n = (int)q.size(), l = 0, r = -1;
 for (int i = 0; i < n; i++) {
```

```
if (i && !argCmp(q[i].dir, q[i-1].dir)) continue;
while (l < r && cov(q[i], q[r-1], q[r])) --r;
while (l < r && cov(q[i], q[l], q[l+1])) ++l;
q[++r] = q[i];
}
while (l < r && cov(q[l], q[r-1], q[r])) --r;
while (l < r && cov(q[r], q[l], q[l+1])) ++l;
n = r - l + 1; // q[l ... r] are the lines
if (n <= 1 || !argCmp(q[l].dir, q[r].dir)) return 0;
vector<PTF> pt(n);
for (int i = 0; i < n; i++)
pt[i] = intersect(q[i+l], q[(i+1)%n+l]);
return area(pt);
} // test @ 2020 Nordic NCPC : BigBrother</pre>
```

6.8 SegmentDist (Sausage) [9d8603]

```
// be careful of abs<complex<int>> (replace _abs below)
llf PointSegDist(P A, Seg B) {
   if (B.dir == P(0)) return _abs(A - B.st);
   if (sgn(dot(A - B.st, B.dir)) *
      sgn(dot(A - B.ed, B.dir)) <= 0)
      return abs(cross(A - B.st, B.dir)) / _abs(B.dir);
   return min(_abs(A - B.st), _abs(A - B.ed));
}
llf SegSegDist(const Seg &s1, const Seg &s2) {
   if (isInter(s1, s2)) return 0;
   return min({
      PointSegDist(s1.st, s2),
      PointSegDist(s2.st, s1),
      PointSegDist(s2.ed, s1) });
} // test @ QOJ2444 / PTZ19 Summer.D3</pre>
```

6.9 Rotating Sweep Line [1d9b4d]

```
void rotatingSweepLine(P a[], int n) {
 vector<pair<int,int>> ls; ls.reserve(n*(n-1)/2);
 for (int i = 0; i < n; ++i)</pre>
  for (int j = i + 1; j < n; ++j)
   ls.emplace_back(i, j);
 sort(all(ls), [&a](auto &u, auto &v){
  P zu = a[u.first] - a[u.second];
P zv = a[v.first] - a[v.second];
  int s = sgn(RE(zu)) * sgn(RE(zv));
  if (s == 0) return RE(zu) != 0;
  return sgn(cross(zu, zv)) * s > 0;
 });
 vector<int> idx(n), p(n);
 iota(all(idx), 0);
 sort(all(idx), [&a](int i, int j) {
  return cmpxy(a[i], a[j]); });
 for (int i = 0; i < n; ++i) p[idx[i]] = i;</pre>
 for (auto [i, j]: ls) {
  // do here
  assert (abs(p[i] - p[j]) == 1);
  swap(p[i], p[j]); idx[p[i]] = i; idx[p[j]] = j;
 } // consider swap same slope together?
```

6.10 Point In Simple Polygon [22ef0b]

```
bool PIP(vector<P> &p, P z, bool strict = true) {
  int cnt = 0, n = p.size();
  for (int i = 0; i < n; i++) {
    P A = p[i], B = p[(i + 1) % n];
    if (isInter(Seg(A, B), z)) return !strict;
    auto zy = IM(z), Ay = IM(A), By = IM(B);
    cnt ^= ((zy<Ay) - (zy<By)) * ori(z, A, B) > 0;
  }
  return cnt;
}
```

6.11 Point In Hull (Fast) [906873]

```
bool PIH(vector<P> &h, P z, bool strict = true) {
  int n = (int)h.size(), a = 1, b = n - 1, r = !strict;
  if (n < 3) return r && isInter(Seg(h[0], h[n-1]), z);
  if (ori(h[0],h[a],h[b]) > 0) swap(a, b);
  if (ori(h[0],h[a],z) >= r || ori(h[0],h[b],z) <= -r)
  return false;
  while (abs(a - b) > 1) {
  int c = (a + b) / 2;
  (ori(h[0], h[c], z) > 0 ? b : a) = c;
  }
```

```
return ori(h[a], h[b], z) < r;
}</pre>
```

6.12 Tangent of Points To Hull [6d7cd7]

```
pair<int, int> get_tangent(const vector<P> &v, P p) {
   const auto gao = [&, N = int(v.size())](int s) {
      const auto lt = [&](int x, int y) {
        return ori(p, v[x % N], v[y % N]) == s; };
      int l = 0, r = N; bool up = lt(0, 1);
      while (r - l > 1) {
        int m = (l + r) / 2;
        if (lt(m, 0) ? up : !lt(m, m+1)) r = m;
        else l = m;
    }
    return (lt(l, r) ? r : l) % N;
}; // test @ codeforces.com/gym/101201/problem/E
    return {gao(-1), gao(1)}; // (a,b):ori(p,v[a],v[b])<0}
} // plz ensure that point strictly out of hull</pre>
```

6.13 Circle Class & Intersection [5111af]

```
llf FMOD(llf x) {
 if (x < -PI) x += PI * 2;
 if (x > PI) x -= PI * 2;
 return x;
struct Cir { PTF o; llf r; };
// be carefule when tangent
vector<llf> intersectAngle(Cir a, Cir b) {
PTF dir = b.o - a.o; llf d2 = norm(dir);

if (norm(a.r - b.r) >= d2) { // norm(x) := |x|^2
  if (a.r < b.r) return {-PI, PI}; // a in b</pre>
  else return {}; // b in a
 } else if (norm(a.r + b.r) <= d2) return {};</pre>
 llf dis = abs(dir), theta = arg(dir);
 llf phi = acos((a.r * a.r + d2 - b.r * b.r) /
(2 * a.r * dis)); // is acos_safe needed ?
llf L = FMOD(theta - phi), R = FMOD(theta + phi);
 return { L, R };
vector<PTF> intersectPoint(Cir a, Cir b) {
 llf d = abs(a.o - b.o);
if (d > b.r+a.r || d < abs(b.r-a.r)) return {};
llf dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;</pre>
 PTF dir = (a.o - b.o) / d;
 PTF u = dir * d1 + b.o;
 PTF v = rot90(dir) * sqrt(max(0.0L, b.r*b.r-d1*d1));
 return {u + v, u - v};
} // test @ AOJ CGL probs
```

6.14 Circle Common Tangent [5ff02c]

```
// be careful of tangent / exact same circle
// sign1 = 1 for outer tang, -1 for inter tang
vector<Line> common_tan(const Cir &a, const Cir &b, int
     sign1) {
if (norm(a.o - b.o) < eps) return {};</pre>
llf d = abs(a.o - b.o), c = (a.r - sign1 * b.r) / d;
PTF v = (b.o - a.o) / d;
if (c * c > 1) return {};
if (abs(c * c - 1) < eps) {
 PTF p = a.o + c * v * a.r;
 return {Line(p, p + rot90(b.o - a.o))};
vector<Line> ret; llf h = sqrt(max(0.0L, 1-c*c));
for (int sign2 : {1, -1}) {
 PTF n = c * v + sign2 * h * rot90(v);
 PTF p1 = a.o + n * a.r;
 PTF p2 = b.o + n * (b.r * sign1);
 ret.emplace_back(p1, p2);
return ret;
```

6.15 Line-Circle Intersection [12b42a]

```
vector<PTF> LineCircleInter(PTF p1, PTF p2, PTF o, llf
    r) {
    PTF ft = p1 + project(o-p1, p2-p1), vec = p2-p1;
    llf dis = abs(o - ft);
    if (abs(dis - r) < eps) return {ft};
    if (dis > r) return {};
    vec = vec * sqrt(r * r - dis * dis) / abs(vec);
    return {ft + vec, ft - vec}; // sqrt_safe?
}
```

6.16 Poly-Circle Intersection [242a4e]

```
// Divides into multiple triangle, and sum up
  from 8BQube, test by HDU2892 & AOJ CGL_7_H
llf _area(PTF pa, PTF pb, llf r) {
 if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
 if (abs(pb) < eps) return 0;</pre>
 llf S, h, theta;
 llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
 llf cB = dot(pb, pb-pa) / a / c, B = acos_safe(cB);
 llf cC = dot(pa, pb) / a / b, C = acos_safe(cC);
 if (a > r) {
  S = (C / 2) * r * r; h = a * b * sin(C) / c;
if (h < r && B < PI / 2)
   S = (acos\_safe(h/r)*r*r - h*sqrt\_safe(r*r-h*h));
 } else if (b > r) {
theta = PI - B - asin_safe(sin(B) / r * a);
  S = 0.5 * a*r*sin(theta) + (C-theta)/2 * r * r;
 } else
  S = 0.5 * sin(C) * a * b;
 return S;
llf area_poly_circle(const vector<PTF> &poly, PTF 0,
    llf r) {
 llf S = 0;
 for (int i = 0, N = poly.size(); i < N; ++i)</pre>
  S += _area(poly[i] - 0, poly[(i + 1) % N] - 0, r) *
     ori(0, poly[i], poly[(i + 1) % N]);
 return abs(S);
}
```

6.17 Minimum Covering Circle [3a9017]

```
// be careful of type
Cir getCircum(P a, P b, P c){
 llf a1 = a.x-b.x, b1 = a.y-b.y;
 llf c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
 llf a2 = a.x-c.x, b2 = a.y-c.y;
 llf c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
 Cir cc;
 cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
 cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
 cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
 return cc;
Cir minCircleCover(vector<P> &pts) {
 shuffle(pts.begin(), pts.end(), mt19937(114514));
 Cir c = { pts[0], 0 };
for(int i = 0; i < (int)pts.size(); i++) {
  if (dist(pts[i], c.o) <= c.r) continue;</pre>
  c = { pts[i], 0 };
for (int j = 0; j < i; j++) {</pre>
   if (dist(pts[j], c.o) <= c.r) continue;
c.o = (pts[i] + pts[j]) / llf(2);</pre>
   c.r = dist(pts[i], c.o);
   for (int k = 0; k < j; k++) {</pre>
    if (dist(pts[k], c.o) <= c.r) continue;</pre>
    c = getCircum(pts[i], pts[j], pts[k]);
  }
 return c;
```

6.18 Circle Union [1a5265]

```
#define eb emplace_back
struct Teve { // test@SPOJ N=1000, 0.3~0.5s
PTF p; llf a; int add; // point, ang, add
Teve(PTF x, llf y, int z) : p(x), a(y), add(z) {}
bool operator<(Teve &b) const { return a < b.a; }
};
// strict: x = 0, otherwise x = -1
bool disjunct(Cir &a, Cir &b, int x)
{ return sgn(abs(a.o - b.o) - a.r - b.r) > x; }
bool contain(Cir &a, Cir &b, int x)
{ return sgn(a.r - b.r - abs(a.o - b.o)) > x; }
vector<llf> CircleUnion(vector<Cir> &c) {
// area[i] : area covered by at least i circles
int N = (int)c.size(); vector<llf> area(N + 1);
vector<vector<int>> overlap(N, vector<int>(N));
auto g = overlap; // use simple 2darray to speedup
for (int i = 0; i < N; ++i)
for (int j = 0; j < N; ++j) {</pre>
```

```
/* c[j] is non-strictly in c[i]. */
  overlap[i][j] = i != j &&
   (sgn(c[i].r - c[j].r) > 0 | |
    (sgn(c[i].r - c[j].r) == 0 \&\& i < j)) \&\&
   contain(c[i], c[j], -1);
for (int i = 0; i < N; ++i)</pre>
 for (int j = 0; j < N; ++j)
  g[i][j] = i != j && !(overlap[i][j] ||
overlap[j][i] || disjunct(c[i], c[j], -1));
for (int i = 0; i < N; ++i) {</pre>
 vector<Teve> eve; int cnt = 1;
 for (int j = 0; j < N; ++j) cnt += overlap[j][i];</pre>
  / if (cnt > 1) continue; (if only need area[1])
 for (int j = 0; j < N; ++j) if (g[i][j]) {</pre>
  auto IP = intersectPoint(c[i], c[j]);
  PTF aa = IP[1], bb = IP[0];
  llf A = arg(aa - c[i].o), B = arg(bb - c[i].o);
  eve.eb(bb, B, 1); eve.eb(aa, A, -1);
  if (B > A) ++cnt;
 if (eve.empty()) area[cnt] += PI*c[i].r*c[i].r;
 else {
  sort(eve.begin(), eve.end());
  eve.eb(eve[0]); eve.back().a += PI * 2;
  for (size_t j = 0; j + 1 < eve.size(); j++) {</pre>
   cnt += eve[j].add;
   area[cnt] += cross(eve[j].p, eve[j+1].p) *.5;
   llf t = eve[j + 1].a - eve[j].a;
   area[cnt] += (t-sin(t)) * c[i].r * c[i].r *.5;
 }
return area;
```

6.19 Polygon Union [2bff43]

```
llf rat(P a, P b) { return sgn(RE(b)) ? llf(RE(a))/RE(b
    ) : llf(IM(a))/IM(b); }
llf polyUnion(vector<vector<P>>& poly) {
llf ret = 0; // area of poly[i] must be non-negative
rep(i,0,sz(poly)) rep(v,0,sz(poly[i])) {
 P A = poly[i][v], B = poly[i][(v + 1) % sz(poly[i])];
 vector<pair<llf, int>> segs{{0, 0}, {1, 0}};
 rep(j,0,sz(poly)) if (i != j) {
   rep(u,0,sz(poly[j])) {
   P C = poly[j][u], D = poly[j][(u + 1) % sz(poly[j])
    if (int sc = ori(A, B, C), sd = ori(A, B, D); sc !=
    sd) {
    llf sa = cross(D-C, A-C), sb = cross(D-C, B-C);
    if (min(sc, sd) < 0)
     segs.emplace_back(sa / (sa - sb), sgn(sc - sd));
    } else if (!sc && !sd && j<i && sgn(dot(B-A,D-C))</pre>
    >0){
    segs.emplace_back(rat(C - A, B - A), 1);
     segs.emplace_back(rat(D - A, B - A), -1);
  }
 }
 sort(segs.begin(), segs.end());
  for (auto &s : segs) s.first = clamp<llf>(s.first, 0,
    1);
 llf sum = 0;
 int cnt = segs[0].second;
 rep(j,1,sz(segs)) {
  if (!cnt) sum += segs[j].first - segs[j - 1].first;
  cnt += segs[j].second;
 ret += cross(A,B) * sum;
return ret / 2;
```

6.20 3D Convex Hull [93b153]

```
// return the faces with pt indexes
struct P3 { lld x,y,z;
P3 operator * (const P3 &b) const {
 return(P3){y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
} };
struct Face { int a, b, c;
Face(int ta,int tb,int tc):a(ta),b(tb),c(tc){} };
```

```
P3 ver(P3 a, P3 b, P3 c) { return (b - a) * (c - a); }
// plz ensure that first 4 points are not coplanar
// all points coplanar case will WA
vector<Face> convex_hull_3D(const vector<P3> &pt) {
 int n = int(pt.size()); vector<Face> now;
 if (n <= 2) return {}; // be careful about edge case</pre>
 vector<vector<int>> flag(n, vector<int>(n));
 now.emplace_back(0,1,2); now.emplace_back(2,1,0);
 for (int i = 3; i < n; i++) {</pre>
  vector<Face> next;
  for (const auto &f : now) {
   lld d = (pt[i] - pt[f.a]).dot(
     ver(pt[f.a], pt[f.b], pt[f.c]));
   if (d <= 0) next.push_back(f);</pre>
   int ff = (d > 0) - (d < 0);
   flag[f.a][f.b]=flag[f.c][f.a]=ff;
  for (const auto &f : now) {
   const auto F = [\&](int x, int y) \{
    if (flag[x][y] > 0 && flag[y][x] <= 0)</pre>
     next.emplace_back(x, y, i);
   F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
  7
  now = next;
 return now;
// delaunay: facets with negative z normal of
// convexhull of (x, y, x^2 + y^2)
```

6.21 Delaunay [7f0d57]

```
/* A triangulation such that no points will strictly
inside circumcircle of any triangle.
find(root, p) : return a triangle contain given point
add_point : add a point into triangulation
Region of triangle u: iterate each u.e[i].tri,
each points are u.p[(i+1)\%3], u.p[(i+2)\%3]
Voronoi diagram: for each triangle in `res`
the bisector of all its edges will split the region. */
#define L(i) ((i)==0 ? 2 : (i)-1)
#define R(i) ((i)==2 ? 0 : (i)+1)
#define F3 for (int i = 0; i < 3; i++)
bool in_cc(const array<P,3> &p, P q) {
  i128 det = 0;
  F3 det += i128(norm(p[i]) - norm(q)) *
   cross(p[R(i)] - q, p[L(i)] - q);
  return det > 0;
struct Tri;
struct E {
 Tri *t; int side; E() : t(0), side(0) {}
 E(Tri *t_, int side_) : t(t_), side(side_){}
struct Tri {
 array<P,3> p; array<Tri*,3> ch; array<E,3> e;
 Tri(Pa = 0, Pb = 0, Pc = 0) : p{a, b, c}, ch{} {}
 bool has_chd() const { return ch[0] != nullptr; }
 bool contains(P q) const {
  F3 if (ori(p[i], p[R(i)], q) < 0) return false;
  return true;
} pool[maxn * 10], *it;
void link(E a, E b) {
 if (a.t) a.t->e[a.side] = b;
 if (b.t) b.t->e[b.side] = a;
struct Trigs {
 Tri *root;
 Trigs() { // should at least contain all points
  root = // C = 100*MAXC^2 or just MAXC?
   new(it++) Tri(P(-C, -C), P(C*2, -C), P(-C, C*2));
 void add_point(P p) { add_point(find(p, root), p); }
 static Tri* find(P p, Tri *r) {
  while (r->has_chd()) for (Tri *c: r->ch)
    if (c && c->contains(p)) { r = c; break; }
  return r;
 void add_point(Tri *r, P p) {
  array<Tri*, 3> t; /* split into 3 triangles */
  F3 t[i] = new (it++) Tri(r->p[i], r->p[R(i)], p);
```

```
F3 link(E(t[i], 0), E(t[R(i)], 1));
 F3 link(E(t[i], 2), r->e[L(i)]);
 r->ch = t;
 F3 flip(t[i], 2);
void flip(Tri* A, int a) {
 auto [B, b] = A->e[a]; /* flip edge between A,B */
 if (!B || !in_cc(A->p, B->p[b])) return;
 Tri *X = new(it++)Tri(A->p[R(a)],B->p[b],A->p[a]);
 Tri *Y = new(it++)Tri(B->p[R(b)],A->p[a],B->p[b]);
 link(E(X,0), E(Y,0));
 link(E(X,1), A->e[L(a)]); link(E(X,2), B->e[R(b)]);
 link(E(Y,1), B->e[L(b)]); link(E(Y,2), A->e[R(a)]);
 A \rightarrow ch = B \rightarrow ch = \{X, Y, nullptr\};
  flip(X, 1); flip(X, 2); flip(Y, 1); flip(Y, 2);
vector<Tri*> res; set<Tri*> vis;
void go(Tri *now) { // store all tri into res
if (!vis.insert(now).second) return;
if (!now->has_chd()) return res.push_back(now);
for (Tri *c: now->ch) if (c) go(c);
void build(vector<P> ps) {
it = pool; res.clear(); vis.clear();
shuffle(ps.begin(), ps.end(), mt19937(114514));
Trigs tr; for (P p: ps) tr.add_point(p);
go(tr.root); // use `res` afterwards
```

6.22 kd Tree (Nearest Point) [f87996]

```
struct KDTree {
struct Node {
 int x, y, x1, y1, x2, y2, id, f;
 Node *L, *R;
 } tree[maxn], *root;
lld dis2(int x1, int y1, int x2, int y2) {
 lld dx = x1 - x2, dy = y1 - y2;
 return dx * dx + dy * dy;
static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
void init(vector<pair<int,int>> &ip) {
 const int n = ip.size();
 for (int i = 0; i < n; i++) {</pre>
  tree[i].id = i;
  tree[i].x = ip[i].first;
  tree[i].y = ip[i].second;
 root = build_tree(0, n-1, 0);
Node* build_tree(int L, int R, int d) {
 if (L>R) return nullptr;
 int M = (L+R)/2; tree[M].f = d%2;
 nth_element(tree+L,tree+M,tree+R+1,d%2?cmpy:cmpx);
 tree[M].x1 = tree[M].x2 = tree[M].x;
 tree[M].y1 = tree[M].y2 = tree[M].y;
 tree[M].L = build_tree(L, M-1, d+1);
 tree[M].R = build_tree(M+1, R, d+1);
 for (Node *s: {tree[M].L, tree[M].R}) if (s) {
  tree[M].x1 = min(tree[M].x1, s->x1);
  tree[M].x2 = max(tree[M].x2, s\rightarrow x2);
  tree[M].y1 = min(tree[M].y1, s->y1);
  tree[M].y2 = max(tree[M].y2, s->y2);
 return tree+M;
bool touch(int x, int y, lld d2, Node *r){
 lld d = sqrt(d2)+1;
 return x >= r->x1 - d && x <= r->x2 + d &&
        y >= r->y1 - d \&\& y <= r->y2 + d;
using P = pair<lld, int>;
void dfs(int x, int y, P &mn, Node *r) {
 if (!r || !touch(x, y, mn.first, r)) return;
 mn = min(mn, P(dis2(r->x, r->y, x, y), r->id));
  // search order depends on split dim
 if (r->f == 1 ? y < r->y : x < r->x) {
  dfs(x, y, mn, r\rightarrow L);
  dfs(x, y, mn, r\rightarrow R);
 } else {
  dfs(x, y, mn, r->R);
```

```
dfs(x, y, mn, r->L);
 int query(int x, int y) {
  P mn(INF, -1);
  dfs(x, y, mn, root);
  return mn.second;
} tree;
```

6.23 kd Closest Pair (3D ver.) [84d9eb]

```
llf solve(vector<P> v) {
 shuffle(v.begin(), v.end(), mt19937());
 unordered_map<lld, unordered_map<lld,
  unordered_map<lld, int>>> m;
 llf d = dis(v[0], v[1]);
 auto Idx = [\&d] (llf x) \rightarrow lld {
  return round(x * 2 / d) + 0.1; };
 auto rebuild_m = [&m, &v, &Idx](int k) {
  m.clear();
  for (int i = 0; i < k; ++i)
   m[Idx(v[i].x)][Idx(v[i].y)]
    [Idx(v[i].z)] = i;
 }; rebuild_m(2);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
     kz = Idx(v[i].z); bool found = false;
  for (int dx = -2; dx \le 2; ++dx) {
   const lld nx = dx + kx;
   if (m.find(nx) == m.end()) continue;
   auto& mm = m[nx];
   for (int dy = -2; dy <= 2; ++dy) {
    const lld ny = dy + ky;
    if (mm.find(ny) == mm.end()) continue;
    auto& mmm = mm[ny];
    for (int dz = -2; dz <= 2; ++dz) {
     const lld nz = dz + kz;
     if (mmm.find(nz) == mmm.end()) continue;
     const int p = mmm[nz];
     if (dis(v[p], v[i]) < d) {</pre>
      d = dis(v[p], v[i]);
      found = true;
     }
   }
  if (found) rebuild_m(i + 1);
  else m[kx][ky][kz] = i;
 return d;
```

6.24 Simulated Annealing [4e0fe5]

```
llf anneal() {
 mt19937 rnd_engine(seed);
 uniform_real_distribution<llf> rnd(0, 1);
 const llf dT = 0.001;
   Argument p
 llf S_cur = calc(p), S_best = S_cur;
 for (llf T = 2000; T > EPS; T -= dT) {
  // Modify p to p_prime
const llf S_prime = calc(p_prime);
  const llf delta_c = S_prime - S_cur;
  llf prob = min((llf)1, exp(-delta_c / T));
  if (rnd(rnd_engine) <= prob)</pre>
   S_cur = S_prime, p = p_prime;
  if (S_prime < S_best) // find min</pre>
   S_best = S_prime, p_best = p_prime;
 }
 return S_best;
```

Stringology

7.1 Hash [7afe3e]

```
class Hash {
 private:
  static constexpr int P = 127, Q = 1051762951;
  vector<int> h, p;
 public:
  void init(const string &s){
```

```
h.assign(s.size()+1, 0); p.resize(s.size()+1);

for (size_t i = 0; i < s.size(); ++i)
                                                                 7.3 Ex SAM [a56a7c]
                                                                 struct exSAM {
    h[i + 1] = add(mul(h[i], P), s[i]);
                                                                  int len[maxn * 2], link[maxn * 2]; // maxlen, suflink
   generate(p.begin(), p.end(),[x=1,y=1,this]()
                                                                  int next[maxn * 2][maxc], tot; // [0, tot), root = 0
     mutable{y=x;x=mul(x,P);return y;});
                                                                  int ord[maxn * 2]; // topo. order
                                                                  int cnt[maxn * 2]; // occurence
  int query(int l, int r){ // 1-base (l, r]
                                                                  int newnode() {
   return sub(h[r], mul(h[l], p[r-l]));}
                                                                   fill_n(next[tot], maxc, 0);
                                                                   return len[tot] = cnt[tot] = link[tot] = 0, tot++;
7.2 Suffix Array [2846f0]
                                                                  void init() { tot = 0, newnode(), link[0] = -1; }
namespace sfx {
                                                                  int insertSAM(int last, int c) {
bool _{t[maxn * 2]};
                                                                   int cur = next[last][c];
int hi[maxn], rev[maxn];
                                                                   len[cur] = len[last] + 1;
int _s[maxn * 2], sa[maxn * 2], _c[maxn * 2];
                                                                   int p = link[last];
int x[maxn], _p[maxn], _q[maxn * 2];
// sa[i]: sa[i]-th suffix is the
                                                                   while (p != -1 && !next[p][c])
                                                                    next[p][c] = cur, p = link[p];
// i-th lexigraphically smallest suffix.
                                                                   if (p == -1) return link[cur] = 0, cur;
// hi[i]: longest common prefix
                                                                   int q = next[p][c];
// of suffix sa[i] and suffix sa[i - 1].
                                                                   if (len[p] + 1 == len[q]) return link[cur] = q, cur;
void pre(int *a, int *c, int n, int z) {
                                                                   int clone = newnode();
 memset(a, 0, sizeof(int) * n);
                                                                   for (int i = 0; i < maxc; ++i)</pre>
 memcpy(x, c, sizeof(int) * z);
                                                                    next[clone][i] = len[next[q][i]] ? next[q][i] : 0;
                                                                   len[clone] = len[p] + 1;
void induce(int *a,int *c,int *s,bool *t,int n,int z){
                                                                   while (p != -1 && next[p][c] == q)
 memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
                                                                    next[p][c] = clone, p = link[p];
                                                                   link[link[cur] = clone] = link[q];
  if (a[i] && !t[a[i] - 1])
                                                                   link[q] = clone;
   a[x[s[a[i] - 1]]++] = a[i] - 1;
                                                                   return cur;
 memcpy(x, c, sizeof(int) * z);
 for (int i = n - 1; i >= 0; --i)
                                                                  void insert(const string &s) {
  if (a[i] && t[a[i] - 1])
                                                                   int cur = 0:
   a[--x[s[a[i] - 1]]] = a[i] - 1;
                                                                   for (auto ch : s) {
                                                                    int &nxt = next[cur][int(ch - 'a')];
void sais(int *s, int *a, int *p, int *q,
bool *t, int *c, int n, int z) {
bool uniq = t[n - 1] = true;
                                                                    if (!nxt) nxt = newnode();
                                                                    cnt[cur = nxt] += 1;
 int nn=0, nmxz=-1, *nsa = a+n, *ns=s+n, last=-1;
 memset(c, 0, sizeof(int) * z);
for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
                                                                  void build() {
                                                                   queue<int> q; q.push(0);
 for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
                                                                   while (!q.empty()) {
 if (uniq) {
                                                                    int cur = q.front(); q.pop();
  for (int i = 0; i < n; ++i) a[--c[s[i]]] = i;</pre>
                                                                    for (int i = 0; i < maxc; ++i)</pre>
  return:
                                                                     if (next[cur][i]) q.push(insertSAM(cur, i));
 for (int i = n - 2; i >= 0; --i)
                                                                   vector<int> lc(tot);
  t[i] = (s[i]==s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
                                                                   for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
 pre(a, c, n, z);
for (int i = 1; i <= n - 1; ++i)</pre>
                                                                   partial_sum(all(lc), lc.begin());
                                                                   for (int i = 1; i < tot; ++i) lenSorted[--lc[len[i]]]</pre>
  if (t[i] && !t[i - 1])
   a[--x[s[i]]] = p[q[i] = nn++] = i;
 induce(a, c, s, t, n, z);

for (int i = 0; i < n; ++i) {
                                                                  void solve() {
                                                                   for (int i = tot - 2; i >= 0; --i)
  if (a[i] && t[a[i]] && !t[a[i] - 1]) {
                                                                    cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
  bool neq = last < 0 ||
memcmp(s + a[i], s + last,</pre>
                                                                };
   (p[q[a[i]] + 1] - a[i]) * sizeof(int));
  ns[q[last = a[i]]] = nmxz += neq;
                                                                 7.4 Z value [6a7fd0]
 }}
                                                                 vector<int> Zalgo(const string &s) {
 sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
                                                                  vector<int> z(s.size(), s.size());
 pre(a, c, n, z);

for (int i = nn - 1; i >= 0; --i)
                                                                  for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
                                                                   int j = clamp(r - i, 0, z[i - l]);
  a[--x[s[p[nsa[i]]]]] = p[nsa[i]];
                                                                   for (; i + j < z[0] and s[i + j] == s[j]; ++j);
 induce(a, c, s, t, n, z);
                                                                   if (i + (z[i] = j) > r) r = i + z[l = i];
void build(const string &s) {
                                                                  return z;
 const int n = int(s.size());
 for (int i = 0; i < n; ++i) _s[i] = s[i];</pre>
 _s[n] = 0; // s shouldn't contain 0
                                                                 7.5 Manacher [365720]
 sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
 for(int i = 0; i < n; ++i) rev[sa[i] = sa[i+1]] = i;</pre>
                                                                 int z[maxn];
 int ind = hi[0] = 0;
                                                                 int manacher(const string& s) {
                                                                  string t = ".";
 for (int i = 0; i < n; ++i) {</pre>
  if (!rev[i]) {
                                                                  for(char c: s) t += c, t += '.';
                                                                  int l = 0, r = 0, ans = 0;
for (int i = 1; i < t.length(); ++i) {
    z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
   ind = 0;
   continue;
                                                                   while (i - z[i] >= 0 && i + z[i] < t.length()) {</pre>
  while (i + ind < n &&</pre>
   s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
                                                                    if(t[i - z[i]] == t[i + z[i]]) ++z[i];
  hi[rev[i]] = ind ? ind-- : 0;
                                                                    else break;
                                                                   if (i + z[i] > r) r = i + z[i], l = i;
| }}
```

```
for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);</pre>
return ans;
```

7.6 Lexico Smallest Rotation [0e9fb8]

```
string mcp(string s) {
int n = s.length();
s += s; int i = 0, j = 1;
while (i < n && j < n) {
  int k = 0;
  while (k < n \& s[i + k] == s[j + k]) k++;
  ((s[i + k] \le s[j + k]) ? j : i) += k + 1;
  j += (i == j);
return s.substr(i < n ? i : j, n);</pre>
```

7.7 Main Lorentz [615b8f]

```
vector<pair<int, int>> rep[kN]; // 0-base [l, r]
void main_lorentz(const string &s, int sft = 0) {
 const int n = s.size();
 if (n == 1) return;
 const int nu = n / 2, nv = n - nu;
 const string u = s.substr(0, nu), v = s.substr(nu),
    ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend());
 main_lorentz(u, sft), main_lorentz(v, sft + nu);
 auto get_z = [](const vector<int> &z, int i) {
  return (0 <= i and i < (int)z.size()) ? z[i] : 0; };</pre>
 auto add_rep = [&](bool left, int c, int l, int k1,
    int k2) {
  const int L = max(1, l - k2), R = min(l - left, k1);
  if (L > R) return;
  if (left) rep[l].emplace_back(sft + c - R, sft + c -
  else rep[l].emplace_back(sft + c - R - l + 1, sft + c
     - L - l + 1);
 for (int cntr = 0; cntr < n; cntr++) {</pre>
  int l, k1, k2;
  if (cntr < nu) {</pre>
  l = nu - cntr;
  k1 = get_z(z1, nu - cntr);
  k2 = get_z(z2, nv + 1 + cntr);
  } else {
  l = cntr - nu + 1;
   k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
  k2 = get_z(z4, (cntr - nu) + 1);
  if (k1 + k2 >= 1)
   add_rep(cntr < nu, cntr, l, k1, k2);</pre>
}
```

7.8 BWT [5a9b3a]

```
vector<int> v[SIGMA];
void BWT(char *ori, char *res) {
  // make ori -> ori + ori
// then build suffix array
void iBWT(char *ori, char *res) {
 for (int i = 0; i < SIGMA; i++) v[i].clear();</pre>
 const int len = strlen(ori);
 for (int i = 0; i < len; i++)</pre>
  v[ori[i] - 'a'].push_back(i);
 vector<int> a;
 for (int i = 0, ptr = 0; i < SIGMA; i++)</pre>
  for (int j : v[i]) {
   a.push_back(j);
   ori[ptr++] = 'a' + i;
 for (int i = 0, ptr = 0; i < len; i++) {</pre>
 res[i] = ori[a[ptr]];
  ptr = a[ptr];
 res[len] = 0;
```

7.9 Palindromic Tree [0673ee]

```
struct PalindromicTree {
 struct node {
  int nxt[26], f, len; // num = depth of fail link
                  // = #pal_suffix of this node
  int cnt, num;
  node(int l = 0) : nxt{}, f(0), len(l), cnt(0), num(0)
 };
 vector<node> st; vector<char> s; int last, n;
 void init() {
  st.clear(); s.clear();
  last = 1; n = 0;
  st.push_back(0); st.push_back(-1);
  st[0].f = 1; s.push_back(-1);
 int getFail(int x) {
  while (s[n - st[x].len - 1] != s[n]) x = st[x].f;
  return x;
 void add(int c) {
  s.push_back(c -= 'a'); ++n;
  int cur = getFail(last);
  if (!st[cur].nxt[c]) {
   int now = st.size();
   st.push_back(st[cur].len + 2);
   st[now].f = st[getFail(st[cur].f)].nxt[c];
   st[cur].nxt[c] = now;
   st[now].num = st[st[now].f].num + 1;
  last = st[cur].nxt[c]; ++st[last].cnt;
 void dpcnt() { // cnt = #occurence in whole str
  for (int i = st.size() - 1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
 int size() { return st.size() - 2; }
} pt;
/* usage
string s; cin >> s; pt.init();
for (int i = 0; i < SZ(s); i++) {
 int prvsz = pt.size(); pt.add(s[i]);
 if (prvsz != pt.size()) {
 int r = i, l = r - pt.st[pt.last].len + 1;
  // pal @ [l,r]: s.substr(l, r-l+1)
} */
```

Misc 8

8.1 Theorems

Sherman-Morrison formula

$$(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1+v^{\mathsf{T}}A^{-1}u}$$

Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} =$ d(i), $L_{ij} = -c$ where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

Cayley's Formula

- Given a degree sequence d_1, d_2, \ldots, d_n for each labeled vertices,
- there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees. Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1, 2, \dots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.

Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \le k \le n$.

Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

Euler's planar graph formula

```
V - E + F = C + 1. E \le 3V - 6 (when V \ge 3)
```

Pick's theorem

For simple polygon, when points are all integer, we have A #{lattice points in the interior} + $\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

Matroid Intersection

iteration of Bellman-Ford is 2R+1

Given matroids $M_1=(G,I_1),M_2=(G,I_2),$ find maximum $S\in I_1\cap I_2.$ For each iteration, build the directed graph and find a shortest path from s to t. $\cdot s\to x:S\sqcup\{x\}\in I_1$ $\cdot x\to t:S\sqcup\{x\}\in I_2$ $\cdot y\to x:S\setminus\{y\}\sqcup\{x\}\in I_2$ $\cdot y\to x:S\setminus\{y\}\sqcup\{x\}\in I_2$ (y is in the unique circuit of $S\sqcup\{x\}$) $\cdot x\to y:S\setminus\{y\}\sqcup\{x\}\in I_2$ (y is in the unique circuit of $S\sqcup\{x\}$) Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|.$ In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Use Bellman-Ford to find the weighted shortest path. The maximum

8.2 Weight Matroid Intersection [c376a9]

```
struct Matroid {
Matroid(bitset<N>); // init from an independent set
bool can_add(int); // check if break independence
Matroid remove(int); // removing from the set
auto matroid_intersection(const vector<int> &w) {
const int n = w.size(); bitset<N> S;
 for (int sz = 1; sz <= n; sz++) {</pre>
 Matroid M1(S), M2(S); vector<vector<pii>>> e(n + 2);
  for (int j = 0; j < n; j++) if (!S[j]) {</pre>
   if (M1.can_add(j)) e[n].eb(j, -w[j]);
   if (M2.can_add(j)) e[j].eb(n + 1, 0);
 for (int i = 0; i < n; i++) if (S[i]) {</pre>
  Matroid T1 = M1.remove(i), T2 = M2.remove(i);
   for (int j = 0; j < n; j++) if (!S[j]) {</pre>
   if (T1.can_add(j)) e[i].eb(j, -w[j]);
    if (T2.can_add(j)) e[j].eb(i, w[i]);
  }
  } // maybe implicit build graph for more speed
  vector<pii> d(n + 2, \{INF, 0\}); d[n] = \{0, 0\};
 vector<int> prv(n + 2, -1);
  // change to SPFA for more speed, if necessary
 bool upd = 1;
 while (upd) {
   upd = 0;
   for (int u = 0; u < n + 2; u++)
   for (auto [v, c] : e[u]) {
     pii x(d[u].first + c, d[u].second + 1);
     if (x < d[v]) d[v] = x, prv[v] = u, upd = 1;
   }
  if (d[n + 1].first >= INF) break;
 for (int x = prv[n+1]; x!=n; x = prv[x]) S.flip(x);
  // S is the max-weighted independent set w/ size sz
return S;
} // from Nacl
```

8.3 Bitset LCS [5e6c56]

```
scanf("%d%d", &n, &m), u = n / 64 + 1;
for (int i = 1, c; i <= n; i++)
scanf("%d", &c), p[c].set(i);
for (int i = 1, c; i <= m; i++) {
    scanf("%d", &c), (g = f) |= p[c];
    f.shiftLeftByOne(), f.set(0);
    ((f = g - f) ^= g) &= g;
}
printf("%d\n", f.count());</pre>
```

8.4 Prefix Substring LCS [78a378]

```
void all_lcs(string s, string t) { // 0-base
vector<int> h(SZ(t));
iota(ALL(h), 0);
for (int a = 0; a < SZ(s); ++a) {
  int v = -1;
  for (int c = 0; c < SZ(t); ++c)
   if (s[a] == t[c] || h[c] < v)</pre>
```

```
swap(h[c], v);
// LCS(s[0, a], t[b, c]) =
// c - b + 1 - sum([h[i] >= b] | i <= c)
// h[i] might become -1 !!
}
}</pre>
```

8.5 Convex 1D/1D DP [6e0124]

```
struct segment {
 int i, l, r;
 segment() {}
 segment(int a, int b, int c): i(a), l(b), r(c) {}
 auto f = [](int l, int r){return dp[l] + w(l+1, r);}
 dp[0] = 0;
 deque<segment> dq; dq.push_back(segment(0, 1, n));
 for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i);
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
   f(i, dq.back().l)<f(dq.back().i, dq.back().l))</pre>
    dq.pop_back();
  if (dq.size()) {
   int d = 1 << 20, c = dq.back().l;</pre>
   while (d >>= 1) if (c + d <= dq.back().r)</pre>
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
  if (seg.l <= n) dq.push_back(seg);</pre>
```

8.6 ConvexHull Optimization [25eb56]

```
struct L {
 mutable lld a, b, p;
 bool operator<(const L &r) const {</pre>
  return a < r.a; /* here */
 bool operator<(lld x) const { return p < x; }</pre>
lld Div(lld a, lld b) {
 return a / b - ((a ^ b) < 0 && a % b); };</pre>
struct DynamicHull : multiset<L, less<>>> {
 static const lld kInf = 1e18;
 bool Isect(iterator x, iterator y) {
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a)
   x->p = x->b > y->b ? kInf : -kInf; /* here */
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(lld a, lld b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
  while ((y = x) != begin() && (--x)->p >= y->p)
   Isect(x, erase(y));
 lld Query(lld x) { // default chmax
  auto l = *lower_bound(x); // to chmin:
return l.a * x + l.b; // modify the 2 "<>"
};
```

8.7 Josephus Problem [f4494f]

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

8.8 Tree Knapsack [87db92]

```
int dp[N][K]; PII obj[N];
vector<int> G[N];
void dfs(int u, int mx){
    for(int s: G[u]) {
        if(mx < obj[s].first) continue;
        for(int i=0;i<=mx-obj[s].FF;i++)
            dp[s][i] = dp[u][i];
        dfs(s, mx - obj[s].first);
        for(int i=obj[s].FF;i<=mx;i++)
            dp[u][i] = max(dp[u][i],
            dp[s][i - obj[s].FF] + obj[s].SS);
    }
}</pre>
```

8.9 N Queens Problem [31f83e]

```
void solve(VI &ret, int n) { // no sol when n=2,3
  if (n % 6 == 2) {
    for (int i = 2; i <= n; i += 2) ret.push_back(i);
    ret.push_back(3); ret.push_back(1);
    for (int i = 7; i <= n; i += 2) ret.push_back(i);
    ret.push_back(5);
} else if (n % 6 == 3) {
    for (int i = 4; i <= n; i += 2) ret.push_back(i);
    ret.push_back(2);
    for (int i = 5; i <= n; i += 2) ret.push_back(i);
    ret.push_back(1); ret.push_back(3);
} else {
    for (int i = 2; i <= n; i += 2) ret.push_back(i);
    for (int i = 1; i <= n; i += 2) ret.push_back(i);
    for (int i = 1; i <= n; i += 2) ret.push_back(i);
}
</pre>
```

8.10 Stable Marriage

```
l: Initialize m \in M and w \in W to free
2: while \exists free man m who has a woman w to propose to do
3: w \leftarrow first woman on m's list to whom m has not yet proposed
4: if \exists some pair (m', w) then
5: if w prefers m to m' then
6: m' \leftarrow free
7: (m, w) \leftarrow engaged
8: end if
9: else
10: (m, w) \leftarrow engaged
11: end if
12: end while
```

8.11 Binary Search On Fraction [765c5a]

```
struct Q {
ll p, q;
Q go(Q b, ll d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
assert(pred(hi));
bool dir = 1, L = 1, H = 1;
for (; L || H; dir = !dir) {
 ll len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
  if (Q mid = hi.go(lo, len + step);
     mid.p > N || mid.q > N || dir ^ pred(mid))
  else len += step;
 swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
return dir ? hi : lo;
```