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```

#### Basic 1.1 vimrc

```
inoremap {<CR> {<CR>}<ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -g -std=gnu++20 -
    DCKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
    Wconversion -fsanitize=address,undefined,float-
    divide-by-zero,float-cast-overflow && echo success<
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -O2 -g -std=gnu
    ++20 && echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space
    :]' \| md5sum \| cut -c-6
let c_no_curly_error=1
" setxkbmap -option caps:ctrl_modifier
1.2 Debug Macro [a45c59]
#define all(x) begin(x), end(x)
#ifdef CKISEKI
#include <experimental/iterator>
#define safe cerr<<__PRETTY_FUNCTION__<<" line "<</pre>
     __LINE__<<" safe\n"
#define debug(a...) debug_(#a, a)
#define orange(a...) orange_(#a, a)
void debug_(auto s, auto ...a) {
  cerr << "\e[1;32m(" << s << ") = (";</pre>
 int f = 0;
 (..., (cerr << (f++ ? ", " : "") << a));
 cerr << ")\e[0m\n";</pre>
void orange_(auto s, auto L, auto R) {
  cerr << "\e[1;33m[ " << s << " ] = [ ";</pre>
 using namespace experimental;
 copy(L, R, make_ostream_joiner(cerr, ", "));
 cerr << " ]\e[0m\n";</pre>
}
#else
#define safe ((void)0)
#define debug(...) safe
#define orange(...) safe
#endif
1.3 SVG Writer [85759e]
#ifdef CKISEKI
class SVG {
 void p(string_view s) { o << s; }</pre>
 void p(string_view s, auto v, auto... vs) {
  auto i = s.find('$');
  o << s.substr(0, i) << v, p(s.substr(i + 1), vs...);
 ofstream o; string c = "red";
public:
 SVG(auto f,auto x1,auto y1,auto x2,auto y2) : o(f) {
  p("<svg xmlns='http://www.w3.org/2000/svg'
   "viewBox='$ $ $'>\n"
   "<style>*{stroke-width:0.5%;}</style>\n",
   x1, -y2, x2 - x1, y2 - y1);}
 ~SVG() { p("</svg>\n"); }
 void color(string nc) { c = nc; }
 void line(auto x1, auto y1, auto x2, auto y2) {
p("<line x1='$' y1='$' x2='$' y2='$' stroke='$'/>\n",
   x1, -y1, x2, -y2, c); }
 void circle(auto x, auto y, auto r) {
p("<circle cx='$' cy='$' r='$' stroke='$' "</pre>
   "fill='none'/>\n", x, -y, r, c); }
 void text(auto x, auto y, string s, int w = 12) {
  p("<text x='$' y='$' font-size='$px'>$</text>\n",
   x, -y, w, s); }
}; // write wrapper for complex if use complex
#else
struct SVG { SVG(auto ...) {} }; // you know how to
#endif
1.4 Pragma Optimization [6006f6]
\verb"pragma" GCC" optimize("Ofast, no-stack-protector")
```

#pragma GCC optimize("no-math-errno,unroll-loops")

#pragma GCC target("popcnt,abm,mmx,avx,arch=skylake")

\_builtin\_ia32\_ldmxcsr(\_\_builtin\_ia32\_stmxcsr()|0x8040)

#pragma GCC target("sse,sse2,sse3,sse3,sse4")

1.5 IO Optimization [c9494b]

static inline int gc() {

se is nu ru et tgc sc hls cin cino+=j1 sw=2 sts=2 bs=2

mouse=a "encoding=utf-8 ls=2

syn on | colo desert | filetype indent on

int find\_root(int u) {

```
constexpr int B = 1<<20; static char buf[B], *p, *q;</pre>
                                                               int la = 0;
 if (p == q) q = (p = buf) + fread(buf, 1, B, stdin);
                                                               for (access(u); u; u = lc) down(la = u);
return q == buf ? EOF : *p++;
                                                               return la;
                                                              void split(int x, int y) { chroot(x); access(y); }
2
     Data Structure
                                                              void chroot(int u) { access(u); set_rev(u); }
   Dark Magic [095f25]
                                                              /* SPLIT_HASH_HERE */
#include <ext/pb_ds/assoc_container.hpp>
                                                             public:
#include <ext/pb_ds/priority_queue.hpp>
                                                              LCT(int n = 0) : o(n + 1) {}
using namespace __gnu_pbds;
                                                              void set_val(int u, const Val &v) {
  splay(++u); cur.v = v; up(u); }
// heap tags: paring/binary/binomial/rc_binomial/thin
template<typename T>
                                                              void set_sval(int u, const SVal &v) {
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>, \
                                                               access(++u); cur.sv = v; up(u); }
                  pairing_heap_tag>;
                                                              Val query(int x, int y) {
  ′pbds_heap::point_iterator
                                                               split(++x, ++y); return o[y].prod; }
// x = pq.push(10); pq.modify(x, 87); a.join(b);
                                                              SVal subtree(int p, int u) {
// tree tags: rb_tree_tag/ov_tree_tag/splay_tree_tag
                                                               chroot(++p); access(++u); return cur.vir + cur.sv; }
template<typename T>
                                                              bool connected(int u, int v) {
using ordered_set = tree<T, null_type, less<T>,
                                                               return find_root(++u) == find_root(++v); }
   rb_tree_tag, tree_order_statistics_node_update>;
                                                              void link(int x, int y) {
  find_by_order, order_of_key
                                                               chroot(++x); access(++y);
// hash tables: cc_hash_table/gp_hash_table
                                                               o[y].vir = o[y].vir + o[x].sub; up(o[x].pa = y);
2.2 Link-Cut Tree [2aaa19] - 0d97f7/f05d4f/642331
                                                              void cut(int x, int y) {
template <typename Val, typename SVal> class LCT {
                                                               split(++x, ++y); o[y].ch[0] = o[x].pa = 0; up(y); }
struct node {
 int pa, ch[2]; bool rev;
                                                             #undef cur
                                                             #undef lc
 Val v, prod, rprod; SVal sv, sub, vir;
 node() : pa{0}, ch{0, 0}, rev{false}, v{},
                                                             #undef rc
  prod{}, rprod{}, sv{}, sub{}, vir{} {}
                                                                   LiChao Segtree [8e1eaf]
                                                             // cmp(l, r, i) := is l better than r at i?
template <typename L, typename Cmp > class LiChao {
#define cur o[u]
#define lc cur.ch[0]
                                                              int n; vector<L> T; Cmp cmp;
void insert(int l, int r, int o, L ln) {
#define rc cur.ch[1]
vector<node> o;
bool is_root(int u) const {
                                                                // if (ln is empty line) return; // constant
  return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u; }
                                                               int m = (l + r) >> 1;
bool is_rch(int u) const {
                                                               bool atL = cmp(ln, T[o], l);
  return o[cur.pa].ch[1] == u && !is_root(u); }
                                                               if (cmp(ln, T[o], m)) atL ^= 1, swap(T[o], ln);
void down(int u) {
                                                               if (r - l == 1) return;
                                                               if (atL) insert(l, m, o << 1, ln);</pre>
 if (not cur.rev) return;
 for (int c : {lc, rc}) if (c) set_rev(c);
                                                               else insert(m, r, o << 1 | 1, ln);
 cur.rev = false;
                                                              L query(int x, int l, int r, int o) {
void up(int u) {
                                                               if (r - l == 1) return T[o];
                                                               int m = (l + r) >> 1;
 cur.prod = o[lc].prod * cur.v * o[rc].prod;
  cur.rprod = o[rc].rprod * cur.v * o[lc].rprod;
                                                               L s = (x < m ? query(x, l, m, o << 1)
                                                                 : query(x, m, r, o << 1 | 1));
 cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
                                                               return cmp(s, T[o], x) ? s : T[o];
void set_rev(int u) {
 swap(lc, rc), swap(cur.prod, cur.rprod);
cur.rev ^= 1;
                                                             public:
                                                              LiChao(int n_, L init, Cmp &&c) : n(n_{-}), T(n * 4, init)
}
                                                                 ), cmp(c) {}
 /* SPLIT_HASH_HERE */
                                                              void insert(L ln) { insert(0, n, 1, ln); }
void rotate(int u) {
                                                              L query(int x) { return query(x, 0, n, 1); }
 int f = cur.pa, g = o[f].pa, l = is_rch(u);
if (cur.ch[l ^ 1]) o[cur.ch[l ^ 1]].pa = f;
                                                             };
// struct Line { lld a, b; };
  if (not is_root(f)) o[g].ch[is_rch(f)] = u;
                                                             // LiChao lct(
 o[f].ch[l] = cur.ch[l ^ 1], cur.ch[l ^ 1] = f;
                                                                 int(xs.size()), Line{0, INF},
  cur.pa = g, o[f].pa = u; up(f);
                                                                 [&u](const Line &l, const Line &r, int i) {
                                                                  lld x = xs[i];
                                                                  return l.a * x + l.b < r.a * x + r.b;
void splay(int u) {
                                                                 Treap* [ae576c]
 vector<int> stk = {u};
 while (not is_root(stk.back()))
                                                              _gnu_cxx::sfmt19937 rnd(7122); // <ext/random>
   stk.push_back(o[stk.back()].pa);
                                                             namespace Treap {
 while (not stk.empty())
  down(stk.back()), stk.pop_back();
                                                             struct node {
                                                              int size, pri; node *lc, *rc, *pa;
node() : size(1), pri(rnd()), lc(0), rc(0), pa(0) {}
  for (int f = cur.pa; not is_root(u); f = cur.pa) {
   if (!is_root(f))
   rotate(is_rch(u) == is_rch(f) ? f : u);
                                                              void pull() {
  rotate(u);
                                                               size = 1; pa = 0;
                                                               if (lc) { size += lc->size; lc->pa = this;
                                                               if (rc) { size += rc->size; rc->pa = this; }
 up(u);
void access(int x) {
                                                             };
  for (int u = x, last = 0; u; u = cur.pa) {
                                                             int SZ(node *x) { return x ? x->size : 0; }
  splay(u);
                                                             node *merge(node *L, node *R) {
                                                              if (not L or not R) return L ? L : R;
   cur.vir = cur.vir + o[rc].sub - o[last].sub;
                                                              if (L->pri > R->pri)
   rc = last; up(last = u);
 }
                                                               return L->rc = merge(L->rc, R), L->pull(), L;
  splay(x);
                                                              else
                                                               return R->lc = merge(L, R->lc), R->pull(), R;
```

```
void splitBySize(node *o, int k, node *&L, node *&R) {
                                                             set<pii>::iterator addInterval(set<pii>& is, int L, int
 if (not o) L = R = 0;
                                                                  R) {
 else if (int s = SZ(o->lc) + 1; s <= k)
                                                              if (L == R) return is.end();
                                                              auto it = is.lower_bound({L, R}), before = it;
  L=o, splitBySize(o->rc, k-s, L->rc, R), L->pull();
                                                              while (it != is.end() && it->first <= R) {</pre>
 else
  R=o, splitBySize(o->lc, k, L, R->lc), R->pull();
                                                               R = max(R, it->second);
                                                               before = it = is.erase(it);
 // SZ(L) == k
int getRank(node *o) { // 1-base
 int r = SZ(o->lc) + 1;
                                                              if (it != is.begin() && (--it)->second >= L) {
                                                               L = min(L, it->first);
R = max(R, it->second);
 for (; o->pa; o = o->pa)
 if (o->pa->rc == o) r += SZ(o->pa->lc) + 1;
 return r;
                                                               is.erase(it);
  // namespace Treap
                                                              return is.insert(before, {L,R});
2.5 Linear Basis* [138d5d]
template <int BITS, typename S = int> struct Basis {
                                                             void removeInterval(set<pii>% is, int L, int R) {
                                                              if (L == R) return;
 static constexpr S MIN = numeric_limits<S>::min();
                                                              auto it = addInterval(is, L, R);
 array<pair<llu, S>, BITS> b;
                                                              auto r2 = it->second;
 Basis() { b.fill({0, MIN}); }
                                                              if (it->first == L) is.erase(it);
 void add(llu x, S p) {
                                                              else (int&)it->second = L;
  for (int i = BITS-1; i>=0; i--) if (x >> i & 1) {
   if (b[i].first == 0) return b[i]={x, p}, void();
                                                              if (R != r2) is.emplace(R, r2);
   if (b[i].second < p)</pre>
                                                                  Graph
    swap(b[i].first, x), swap(b[i].second, p);
   x ^= b[i].first;
                                                             3.1 SCC [16c7d6]
  }
                                                             class SCC { // test @ library checker
 }
                                                             protected:
 optional<llu> query_kth(llu v, llu k) {
                                                              int n, dfc, nscc; vector<vector<int>> G;
 vector<pair<llu, int>> o;
                                                              vector<int> vis, low, idx, stk;
  for (int i = 0; i < BITS; i++)</pre>
                                                              void dfs(int i) {
   if (b[i].first) o.emplace_back(b[i].first, i);
                                                               vis[i] = low[i] = ++dfc; stk.push_back(i);
  if (k >= (1ULL << o.size())) return {};</pre>
                                                               for (int j : G[i])
  for (int i = int(o.size()) - 1; i >= 0; i--)
                                                                if (!vis[j])
                                                                dfs(j), low[i] = min(low[i], low[j]);
else if (vis[j] != -1)
low[i] = min(low[i], vis[j]);
   if ((k >> i & 1) ^ (v >> o[i].second & 1))
    v ^= o[i].first;
  return v;
                                                               if (low[i] == vis[i])
 Basis filter(S l) {
                                                                 for (idx[i] = nscc++; vis[i] != -1;) {
  Basis res = *this;
                                                                 int x = stk.back(); stk.pop_back();
  for (int i = 0; i < BITS; i++)</pre>
                                                                  idx[x] = idx[i]; vis[x] = -1;
   if (res.b[i].second < l) res.b[i] = {0, MIN};</pre>
  return res;
 }
                                                             public:
                                                              SCC(int n_{-}) : n(n_{-}), dfc(0), nscc(0), G(n),
2.6 Binary Search on Segtree [6c61c0]
                                                               vis(n), low(n), idx(n) {}
// find_first = l -> minimal x s.t. check( [l, x) )
                                                              void add_edge(int u, int v) { G[u].push_back(v); }
// find_last = r \rightarrow maximal x s.t. check([x, r))
                                                              void solve() {
int find_first(int l, auto &&check) {
                                                               for (int i = 0; i < n; i++) if (!vis[i]) dfs(i); }</pre>
 if (l >= n) return n + 1;
                                                              int get_id(int x) { return idx[x]; }
 l += sz; push(l); Monoid sum; // identity
                                                              int count() { return nscc; }
                                                             }; // dag edges point from idx large to idx small
  while ((l & 1) == 0) l >>= 1;
                                                              3.2 2-SAT [ca961f]
  if (auto s = sum + nd[l]; check(s)) {
                                                             struct TwoSat : SCC {
   while (l < sz) {</pre>
                                                              void orr(int x, int y) {
    prop(l); l = (l << 1);
                                                               if ((x ^ y) == 1) return;
    if (auto nxt = sum + nd[l]; not check(nxt))
                                                               add_edge(x ^ 1, y); add_edge(y ^ 1, x);
     sum = nxt, l++;
                                                              vector<int> solve2sat() {
   return l + 1 - sz;
                                                               solve(); vector<int> res(n);
for (int i = 0; i < n; i += 2)
  if (idx[i] == idx[i + 1]) return {};</pre>
  } else sum = s, l++;
 } while (lowbit(l) != l);
 return n + 1;
                                                               for (int i = 0; i < n; i++)</pre>
                                                                res[i] = idx[i] < idx[i ^ 1];
int find_last(int r, auto &&check) {
                                                               return res;
 if (r <= 0) return -1;
                                                              }
 r += sz; push(r - 1); Monoid sum; // identity
                                                             3.3 BCC [6ac6db]
                                                             class BCC {
  while (r > 1 and (r & 1)) r >>= 1;
                                                              int n, ecnt, bcnt;
  if (auto s = nd[r] + sum; check(s)) {
                                                              vector<vector<pair<int, int>>> g;
  while (r < sz) {</pre>
                                                              vector<int> dfn, low, bcc, stk;
    prop(r); r = (r << 1) | 1;
                                                              vector<bool> ap, bridge;
    if (auto nxt = nd[r] + sum; not check(nxt))
                                                              void dfs(int u, int f) {
     sum = nxt, r--;
                                                               dfn[u] = low[u] = dfn[f] + 1;
                                                               int ch = 0;
   return r - sz;
                                                               for (auto [v, t] : g[u]) if (bcc[t] == -1) {
  } else sum = s;
                                                                bcc[t] = 0; stk.push_back(t);
 } while (lowbit(r) != r);
                                                                if (dfn[v]) {
 return -1;
                                                                 low[u] = min(low[u], dfn[v]);
                                                                  continue;
     Interval Container* [edce47]
```

```
++ch, dfs(v, u);
   low[u] = min(low[u], low[v]);
if (low[v] > dfn[u]) bridge[t] = true;
                                                                   for (int i = 0; i < n; i++)</pre>
                                                                    if (in[i] == -1) dfs(dfs, i, -1);
                                                                   for (int i = 0; i < n; i++)
if (dsu.anc(i) == i) id[i] = cnt++;</pre>
   if (low[v] < dfn[u]) continue;</pre>
   ap[u] = true;
   while (not stk.empty()) {
                                                                   vector<vector<int>> comps(cnt);
                                                                   for (int i = 0; i < n; i++)</pre>
    int o = stk.back(); stk.pop_back();
    bcc[o] = bcnt;
                                                                    comps[id[dsu.anc(i)]].push_back(i);
    if (o == t) break;
                                                                   return comps;
                                                                  } // test @ yosupo judge
3.6 Bipolar Orientation [b50cd3]
   bcnt += 1;
                                                                  struct BipolarOrientation {
  ap[u] = ap[u] and (ch != 1 or u != f);
                                                                   int n; vector<vector<int>> g;
                                                                   vector<int> vis, low, pa, sgn, ord;
public:
                                                                   BipolarOrientation(int n_) : n(n_),
                                                                    g(n), vis(n), low(n), pa(n, -1), sgn(n) {}
BCC(int n_{-}) : n(n_{-}), ecnt(0), bcnt(0), g(n), dfn(n),
 low(n), stk(), ap(n) {}
void add_edge(int u, int v) {
                                                                   void dfs(int i) {
                                                                    ord.push_back(i); low[i] = vis[i] = int(ord.size());
  g[u].emplace_back(v, ecnt);
                                                                    for (int j : g[i])
                                                                     if (!vis[j])
  g[v].emplace_back(u, ecnt++);
                                                                      pa[j] = i, dfs(j), low[i] = min(low[i], low[j]);
                                                                      else low[i] = min(low[i], vis[j]);
 void solve() {
  bridge.assign(ecnt, false); bcc.assign(ecnt, -1);
for (int i = 0; i < n; ++i) if (!dfn[i]) dfs(i, i);</pre>
                                                                   vector<int> solve(int S, int T) {
                                                                    g[S].insert(g[S].begin(), T); dfs(S);
int bcc_id(int x) const { return bcc[x]; ]
bool is_ap(int x) const { return ap[x]; }
                                                                    vector<int> nxt(n + 1, n), prv = nxt;
                                                                    nxt[S] = T; prv[T] = S; sgn[S] = -1;
 bool is_bridge(int x) const { return bridge[x]; }
                                                                    for (int i : ord) if (i != S && i != T) {
                                                                     int p = pa[i], l = ord[low[i] - 1];
                                                                      if (sgn[l] > 0) // insert after
      Round Square Tree [cf6d74]
                                                                      nxt[i] = nxt[prv[i] = p], nxt[p] = prv[nxt[p]] = i;
struct RST { // be careful about isolate point
 int n; vector<vector<int>> T;
                                                                      else
 RST(auto &G) : n(int(G.size())), T(n) {
                                                                      prv[i] = prv[nxt[i] = p], prv[p] = nxt[prv[p]] = i;
 vector<int> stk, vis(n), low(n);
                                                                      sgn[p] = -sgn[l];
  auto dfs = [&](auto self, int u, int d) -> void {
   low[u] = vis[u] = d; stk.push_back(u);
                                                                    vector<int> v;
   for (int v : G[u]) if (!vis[v]) {
                                                                    for (int x = S; x != n; x = nxt[x]) v.push_back(x);
                                                                    return v;
    self(self, v, d + 1);
                                                                   } // S, T are unique source / unique sink
void add_edge(int a, int b) {
    if (low[v] == vis[u]) {
     int cnt = int(T.size()); T.emplace_back();
                                                                    g[a].emplace_back(b); g[b].emplace_back(a); }; // 存在 ST 雙極定向 iff 連接 (S,T) 後整張圖點雙連通.7 DMST [f4317e]
     for (int x = -1; x != v; stk.pop_back())
      T[cnt].push_back(x = stk.back());
     T[u].push_back(cnt); // T is rooted
    } else low[u] = min(low[u], low[v]);
                                                                  using lld = int64_t;
   } else low[u] = min(low[u], vis[v]);
                                                                  struct E { int s, t; lld w; }; // O-base
                                                                  struct PQ {
  for (int u = 0; u < n; u++)</pre>
                                                                   struct P {
                                                                    lld v; int i;
   if (!vis[u]) dfs(dfs, u, 1);
 \} // T may be forest; after dfs, stk are the roots
                                                                    bool operator>(const P &b) const { return v > b.v; }
   // test @ 2020 Shanghai K
Edge TCC [5a2668]
                                                                   min_heap<P> pq; lld tag;
vector<vector<int>> ETCC(auto &adj) {
                                                                   void push(P p) { p.v -= tag; pq.emplace(p); }
                                                                   P top() { P p = pq.top(); p.v += tag; return p; }
 const int n = static_cast<int>(adj.size());
 vector<int> up(n), low(n), in, out, nx, id;
                                                                   void join(PQ &b) {
 in = out = nx = id = vector<int>(n, -1);
                                                                    if (pq.size() < b.pq.size())</pre>
 int dfc = 0, cnt = 0; Dsu dsu(n);
                                                                     swap(pq, b.pq), swap(tag, b.tag);
 auto merge = [&](int u, int v) {
                                                                    while (!b.pq.empty()) push(b.top()), b.pq.pop();
 dsu.join(u, v); up[u] += up[v]; };
auto dfs = [&](auto self, int u, int p) -> void {
  in[u] = low[u] = dfc++;
                                                                  vector<int> dmst(const vector<E> &e, int n, int root) {
                                                                   vector<PQ> h(n * 2);
for (int i = 0; i < int(e.size()); ++i)</pre>
  for (int v : adj[u]) if (v != u) {
   if (v == p) { p = -1; continue; }
   if (in[v] == -1) {
                                                                    h[e[i].t].push({e[i].w, i});
    self(self, v, u);
if (nx[v] == -1 && up[v] <= 1) {</pre>
                                                                   vector<int> a(n * 2); iota(all(a), 0);
vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
     up[u] += up[v]; low[u] = min(low[u], low[v]);
                                                                   auto o = [\&](auto Y, int X) \rightarrow int {
     continue;
                                                                    return x==a[x] ? x : a[x] = Y(Y, a[x]); };
                                                                   auto S = [&](int i) { return o(o, e[i].s); };
    if (up[v] == 0) v = nx[v];
                                                                   int pc = v[root] = n;
                                                                   for (int i = 0; i < n; ++i) if (v[i] == -1)</pre>
    if (low[u] > low[v])
     low[u] = low[v], swap(nx[u], v);
                                                                    for (int p = i; v[p]<0 || v[p]==i; p = S(r[p])) {</pre>
    for (; v != -1; v = nx[v]) merge(u, v);
                                                                     if (v[p] == i)
                                                                       for (int q = pc++; p != q; p = S(r[p])) {
   } else if (in[v] < in[u]) {</pre>
    low[u] = min(low[u], in[v]); up[u]++;
                                                                        h[p].tag -= h[p].top().v; h[q].join(h[p]);
                                                                        pa[p] = a[p] = q;
    for (int &x = nx[u]; x != -1 &&
  in[x] <= in[v] && in[v] < out[x]; x = nx[x])</pre>
                                                                     while (S(h[p].top().i) == p) h[p].pq.pop();
v[p] = i; r[p] = h[p].top().i;
     merge(u, x);
    up[u]--;
                                                                   vector<int> ans;
   }
                                                                   for (int i = pc - 1; i >= 0; i--) if (v[i] != n) {
  out[u] = dfc;
                                                                    for (int f = e[r[i]].t; f!=-1 && v[f]!=n; f = pa[f])
```

```
v[f] = n;
                                                                  for (int i = 1; i <= N; i++) X[i] = 1;</pre>
                                                                 for (int t = 0; t < E.size(); t++) {</pre>
  ans.push_back(r[i]);
                                                                  auto [u, v] = E[t];
return ans; // default minimize, returns edgeid array
                                                                  int v0 = v, c = X[u], c0 = c, d;
                                                                  vector<pair<int, int>> L; int vst[kN] = {};
     Dominator Tree [ea5b7c]
                                                                  while (!G[u][v0]) {
                                                                    L.emplace_back(v, d = X[v]);
if (!C[v][c]) for (a=L.size()-1;a>=0;a--)
struct Dominator {
vector<vector<int>> g, r, rdom; int tk;
vector<int> dfn, rev, fa, sdom, dom, val, rp;
                                                                      c = color(u, L[a].first, c);
Dominator(int n) : g(n), r(n), rdom(n), tk(0) {
  dfn = rev = fa = sdom = dom =
                                                                    else if (!C[uj[d]) for (a=L.size()-1;a>=0;a--)
color(u, L[a].first, L[a].second);
   val = rp = vector<int>(n, -1); }
                                                                    else if (vst[d]) break;
                                                                    else vst[d] = 1, v = C[u][d];
 void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
 rev[dfn[x] = tk] = x;
                                                                  if (!G[u][v0]) {
  fa[tk] = sdom[tk] = val[tk] = tk; tk++;
                                                                    for (; v; v = flip(v, c, d), swap(c, d));
  for (int u : g[x]) {
                                                                    if (C[u][c0]) { a = int(L.size()) - 1;
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
                                                                    while (--a >= 0 && L[a].second != c);
                                                                     for(;a>=0;a--)color(u,L[a].first,L[a].second);
  r[dfn[u]].push_back(dfn[x]);
 }
                                                                    } else t--;
}
void merge(int x, int y) { fa[x] = y; }
                                                                 }
int find(int x, int c = 0) {
   if (fa[x] == x) return c ? -1 : x;
                                                                 3.10
                                                                       Centroid Decomp.* [670cdd]
  if (int p = find(fa[x], 1); p != -1) {
  if (sdom[val[x]] > sdom[val[fa[x]]])
                                                                 vector<vector<pair<int, int>>> g; // g[u] = \{(v, w)\}
    val[x] = val[fa[x]];
                                                                 vector<int> pa, dep, vis, sz, mx;
   fa[x] = p;
                                                                 vector<vector<int64_t>> Dist;
                                                                 vector<int64_t> Sub, Sub2;
   return c ? p : val[x];
 } else return c ? fa[x] : val[x];
                                                                 vector<int> Cnt, Cnt2;
                                                                 void DfsSz(vector<int> &tmp, int x) {
vector<int> build(int s, int n) {
  // return the father of each node in dominator tree
                                                                  vis[x] = true, sz[x] = 1, mx[x] = 0;
for (auto [u, w] : g[x]) if (not vis[u]) {
 dfs(s); // p[i] = -2 \text{ if i is unreachable from s}
                                                                   DfsSz(tmp, u); sz[x] += sz[u];
 for (int i = tk - 1; i >= 0; --i) {
                                                                   mx[x] = max(mx[x], sz[u]);
   for (int u : r[i])
    sdom[i] = min(sdom[i], sdom[find(u)]);
                                                                  tmp.push_back(x);
   if (i) rdom[sdom[i]].push_back(i);
   for (int u : rdom[i]) {
                                                                 void DfsDist(int x, int64_t D = 0) {
    int p = find(u);
                                                                  Dist[x].push_back(D); vis[x] = true;
    dom[u] = (sdom[p] == i ? i : p);
                                                                  for (auto [u, w] : g[x])
                                                                    if (not vis[u]) DfsDist(u, D + w);
   if (i) merge(i, rp[i]);
                                                                 void DfsCen(int x, int D, int p) {
  vector<int> tmp; DfsSz(tmp, x);
  vector<int> p(n, -2); p[s] = -1;
  for (int i = 1; i < tk; ++i)</pre>
                                                                  int M = int(tmp.size()), C = -1;
   if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
                                                                  for (int u : tmp)
 for (int i = 1; i < tk; ++i)</pre>
                                                                   if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
  p[rev[i]] = rev[dom[i]];
                                                                  for (int u : tmp) vis[u] = false;
                                                                  DfsDist(C);
  return p;
} // test @ yosupo judge
                                                                  for (int u : tmp) vis[u] = false;
                                                                  pa[C] = p, vis[C] = true, dep[C] = D;
for (auto [u, w] : g[C])
};
3.9
      Edge Coloring [029763]
// \max(d_u) + 1 edge coloring, time: O(NM)
                                                                    if (not vis[u]) DfsCen(u, D + 1, C);
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
                                                                public:
void clear(int N) {
for (int i = 0; i <= N; i++)</pre>
                                                                 Centroid(int N) : g(N), pa(N), dep(N),
                                                                  vis(N), sz(N), mx(N), Dist(N),
Sub(N), Sub2(N), Cnt(N), Cnt2(N) {}
 for (int j = 0; j <= N; j++)</pre>
    C[i][j] = G[i][j] = 0;
                                                                 void AddEdge(int u, int v, int w) {
                                                                  g[u].emplace_back(v, w);
void solve(vector<pair<int, int>> &E, int N) {
int X[kN] = {}, a;
                                                                  g[v].emplace_back(u, w);
auto update = [&](int u) {
 for (X[u] = 1; C[u][X[u]]; X[u]++);
                                                                 void Build() { DfsCen(0, 0, -1); }
                                                                 void Mark(int v) {
};
auto color = [&](int u, int v, int c) {
                                                                  int x = v, z = -1;
 int p = G[u][v];
                                                                  for (int i = dep[v]; i >= 0; --i) {
 G[u][v] = G[v][u] = c;
                                                                    Sub[x] += Dist[v][i], Cnt[x]++;
                                                                    if (z != -1)
 C[u][c] = v, C[v][c] = u;
 C[u][p] = C[v][p] = 0;
                                                                     Sub2[z] += Dist[v][i], Cnt2[z]++;
 if (p) X[u] = X[v] = p;
                                                                    x = pa[z = x];
 else update(u), update(v);
                                                                  }
  return p;
                                                                 int64_t Query(int v) {
};
auto flip = [&](int u, int c1, int c2) {
                                                                  int64_t res = 0;
 int p = C[u][c1];
                                                                  int x = v, z = -1;
                                                                  for (int i = dep[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Cnt[x] * Dist[v][i];
 swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
 if (!C[u][c1]) X[u] = c1;
                                                                    if (z != -1)
                                                                    res -= Sub2[z] + 1LL * Cnt2[z] * Dist[v][i];
 if (!C[u][c2]) X[u] = c2;
  return p;
                                                                    x = pa[z = x];
};
```

```
return res;
  // pa, dep are centroid tree attributes
3.11 Lowbit Decomp. [2d7032]
class LBD {
int n, timer, chains;
vector<vector<int>> G;
vector<int> tl, tr, chain, top, dep, pa;
// chains : number of chain
// tl, tr[u] : subtree interval in the seq. of u
// top[i] : top of the chain of vertex i
// chian[u] : chain id of the chain u is on
void predfs(int u, int f) {
 dep[u] = dep[pa[u] = f] + 1;
  for (int v : G[u]) if (v != f) {
   predfs(v, u);
   if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
   chain[u] = chain[v];
  if (chain[u] == 0) chain[u] = ++chains;
void dfschain(int u, int f, int t) {
 tl[u] = timer++; top[u] = t;
  for (int v : G[u])
  if (v != f and chain[v] == chain[u])
   dfschain(v, u, t);
  for (int v : G[u])
   if (v != f and chain[v] != chain[u])
   dfschain(v, u, v);
  tr[u] = timer;
}
public:
LBD(auto &&G_) : n((int)size(G_)),
 timer(0), chains(0), G(G_{-}), tl(n), tr(n),
   chain(n), top(n + 1, -1), dep(n), pa(n)
  { predfs(0, 0); dfschain(0, 0, 0); }
PII get_subtree(int u) { return {tl[u], tr[u]}; }
vector<PII> get_path(int u, int v) {
 vector<PII> res;
 while (top[u] != top[v]) {
   if (dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
  int s = top[u];
  res.emplace_back(tl[s], tl[u] + 1);
  u = pa[s];
 if (dep[u] < dep[v]) swap(u, v);</pre>
  res.emplace_back(tl[v], tl[u] + 1);
  return res;
}; // 記得在資結上對點的修改要改成對其 dfs 序的修改
3.12 Virtual Tree [44f764]
vector<pair<int, int>> build(vector<int> vs, int r) {
vector<pair<int, int>> res;
sort(vs.begin(), vs.end(), [](int i, int j) {
 return dfn[i] < dfn[j]; });</pre>
 vector<int> s = {r};
for (int v : vs) if (v != r) {
  if (int o = lca(v, s.back()); o != s.back()) {
   while (s.size() >= 2) {
   if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
    res.emplace_back(s[s.size() - 2], s.back());
    s.pop_back();
   if (s.back() != o)
    res.emplace_back(o, s.back()), s.back() = o;
  s.push_back(v);
}
for (size_t i = 1; i < s.size(); ++i)</pre>
 res.emplace_back(s[i - 1], s[i]);
return res; // (x, y): x->y
} // 記得建虛樹會多出 `vs` 以外的點
3.13 Tree Hashing [d6a9f9]
vector<int> g[maxn]; llu h[maxn];
llu F(llu z) { // xorshift64star from iwiwi
z ^= z >> 12; z ^= z << 25; z ^= z >> 27;
return z * 2685821657736338717LL;
llu hsah(int u, int f) {
llu r = 127; // bigger?
for (int v : g[u]) if (v != f) r += hsah(v, u);
```

```
return h[u] = F(r);
  // test @ UOJ 763 & yosupo library checker
3.14 Mo's Algo on Tree
dfs u:
 push u
 iterate subtree
 push u
Let P = LCA(u, v) with St(u) \le St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
      Count Cycles [c7e8f2]
// ord = sort by deg decreasing, rk[ord[i]] = i
// D[i] = edge point from rk small to rk big
for (int x : ord) { // c3
 for (int y : D[x]) vis[y] = 1;
 for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
 for (int y : D[x]) vis[y] = 0;
for (int x : ord) { // c4
 for (int y : D[x]) for (int z : adj[y])
  if (rk[z] > rk[x]) c4 += vis[z]++;
 for (int y : D[x]) for (int z : adj[y])
  if (rk[z] > rk[x]) --vis[z];
  // both are O(M*sqrt(M)), test @ 2022 CCPC guangzhou
3.16 Maximal Clique [2da556]
#define iter(u, B) for (size_t u = B._Find_first(); \
  u < n; u = B._Find_next(u))</pre>
// contain a self loop u to u, than u won't in clique
template <size_t maxn> class MaxClique {
private:
 using bits = bitset<maxn>;
 bits popped, G[maxn], ans;
 size_t deg[maxn], deo[maxn], n;
 void sort_by_degree() {
  popped.reset();
  for (size_t i = 0; i < n; ++i) deg[i] = G[i].count();</pre>
  for (size_t i = 0; i < n; ++i) {</pre>
   size_t mi = maxn, id = 0;
   for (size_t j = 0; j < n; ++j)</pre>
    if (!popped[j] and deg[j] < mi) mi = deg[id = j];</pre>
   popped[deo[i] = id] = 1;
iter(u, G[i]) --deg[u];
 void BK(bits R, bits P, bits X) {
  if (R.count() + P.count() <= ans.count()) return;</pre>
  if (not P.count() and not X.count()) {
   if (R.count() > ans.count()) ans = R;
   return:
  /* greedily chosse max degree as pivot
  bits cur = P \mid X; size_t pv = 0, sz = 0;
  iter(u, cur) if (deg[u] > sz) sz = deg[pv = u];
  cur = P \& \sim G[pv] \& \sim R; */ // or simply choose first bits <math>cur = P \& (\sim G[(P \mid X).\_Find\_first()]) \& \sim R;
  iter(u, cur) {
   R[u] = 1; BK(R, P \& G[u], X \& G[u]);
   R[u] = P[u] = 0, X[u] = 1;
 }
public:
 void init(size_t n_) {
  n = n_; ans.reset();
  for (size_t i = 0; i < n; ++i) G[i].reset();</pre>
 void add_edges(int u, bits S) { G[u] = S; }
void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
 int solve() {
  sort_by_degree(); // or simply iota( deo... )
  for (size_t i = 0; i < n; ++i) deg[i] = G[i].count();</pre>
  bits pob, nob = 0; pob.set();
  for (size_t i = n; i < maxn; ++i) pob[i] = 0;</pre>
  for (size_t i = 0; i < n; ++i) {</pre>
   size_t v = deo[i]; bits tmp; tmp[v] = 1;
   BK(tmp, pob & G[v], nob & G[v]);
   pob[v] = 0, nob[v] = 1;
  return static_cast<int>(ans.count());
 }
3.17
       Maximum Clique [aee5d8]
```

```
constexpr size_t kN = 150; using bits = bitset<kN>;
                                                             while (vst[st] != 2) {
                                                              int v = rho.back(); rho.pop_back();
struct MaxClique {
bits G[kN], cs[kN];
                                                              cycle.emplace_back(v); vst[v]++;
 int ans, sol[kN], q, cur[kN], d[kN], n;
void init(int _n) {
                                                             reverse(all(eid)); eid.resize(cycle.size());
 n = _n;
                                                             return mmc;
  for (int i = 0; i < n; ++i) G[i].reset();</pre>
                                                             3.19
                                                                  Eulerian Trail [8a70bf]
void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
                                                            // g[i] = list of (edge.to, edge.id)
void pre_dfs(vector<int> &v, int i, bits mask) {
                                                            auto euler(int N, int M, int S, const auto &g) {
  if (i < 4) {
                                                             vector<int> iter(N), vis(M), vv, ee;
  for (int x : v) d[x] = (int)(G[x] \& mask).count();
                                                             auto dfs = [&](auto self, int i) -> void {
  sort(all(v), [&](int x, int y) {
                                                              while (iter[i] < ssize(g[i])) {</pre>
    return d[x] > d[y]; });
                                                               auto [j, eid] = g[i][iter[i]++];
                                                               if (vis[eid]) continue;
 vector<int> c(v.size());
                                                               vis[eid] = true; self(self, j);
  cs[1].reset(), cs[2].reset();
                                                               vv.push_back(j); ee.push_back(eid);
  int l = max(ans - q + 1, 1), r = 2, tp = 0, k;
  for (int p : v) {
   for (k = 1; (cs[k] & G[p]).any(); ++k);
                                                             dfs(dfs, S); vv.push_back(S);
   if (k >= r) cs[++r].reset();
                                                             reverse(all(vv)); reverse(all(ee));
   cs[k][p] = 1;
                                                             return pair{vv, ee};
// 需要保證傳入的 g, S degree 符合條件;小心孤點奇點
   if (k < l) v[tp++] = p;
                                                            4
                                                                  Flow & Matching
 for (k = l; k < r; ++k)
for (auto p = cs[k]._Find_first();</pre>
                                                            4.1 HopcroftKarp [930040]
                                                            struct HK {
     p < kN; p = cs[k]._Find_next(p))
                                                             vector<int> l, r, a, p; int ans;
    v[tp] = (int)p, c[tp] = k, ++tp;
                                                             HK(int n, int m, auto \&g) : l(n,-1), r(m,-1), ans(0) {
 dfs(v, c, i + 1, mask);
                                                              for (bool match = true; match;) {
                                                               match = false; a.assign(n, -1); p.assign(n, -1);
void dfs(vector<int> &v, vector<int> &c,
                                                                queue<int> q;
   int i, bits mask) {
                                                               for (int i = 0; i < n; i++)
  while (!v.empty()) {
                                                                if (l[i] == -1) q.push(a[i] = p[i] = i);
   int p = v.back(); v.pop_back(); mask[p] = 0;
                                                                // bitset<maxn> nvis, t; nvis.set();
   if (q + c.back() <= ans) return;</pre>
                                                               while (!q.empty()) {
   cur[q++] = p;
                                                                int z, x = q.front(); q.pop();
   vector<int> nr;
                                                                 if (l[a[x]] != -1) continue;
   for (int x : v) if (G[p][x]) nr.push_back(x);
                                                                for (int y : g[x]) { // or iterate t = g[x]&nvis
   if (!nr.empty()) pre_dfs(nr, i, mask & G[p]);
   else if (q > ans) ans = q, copy_n(cur, q, sol);
                                                                  // nvis.reset(y);
                                                                  if (r[y] == -1) {
   c.pop_back(); --q;
                                                                  for (z = y; z != -1;)
 }
                                                                   r[z] = x, swap(l[x], z), x = p[x];
                                                                  match = true; ++ans; break;
int solve() {
                                                                  } else if (p[r[y]] == -1)
 vector<int> v(n); iota(all(v), 0);
                                                                  q.push(z = r[y]), p[z] = x, a[z] = a[x];
  ans = q = 0; pre_dfs(v, 0, bits(string(n, '1')));
  return ans; // sol[0 ~ ans-1]
icliq; // test @ yosupo
3.18 Min Mean Cycle [e23bc0]
                                                              }
                                                             }
// WARNING: TYPE matters
                                                             4.2
struct Edge { int s, t; llf c; };
                                                                   Kuhn Munkres [74bf6d]
llf solve(vector<Edge> &e, int n) {
                                                            struct KM { // maximize, test @ UOJ 80
                                                             int n, l, r; lld ans; // fl and fr are the match
// O(VE), returns inf if no cycle, mmc otherwise
vector<VI> prv(n + 1, VI(n)), prve = prv;
                                                             vector<lld> hl, hr; vector<int> fl, fr, pre, q;
vector < vector < llf >> d(n + 1, vector < llf > (n, inf));
                                                             void bfs(const auto &w, int s) {
d[0] = vector<llf>(n, 0);
                                                              vector<int> vl(n), vr(n); vector<lld> slk(n, INF);
for (int i = 0; i < n; i++) {</pre>
                                                              l = r = 0; vr[q[r++] = s] = true;
 for (int j = 0; j < (int)e.size(); j++) {</pre>
                                                              auto check = [\&](int x) \rightarrow bool {
   auto [s, t, c] = e[j];
                                                               if (vl[x] || slk[x] > 0) return true;
                                                               vl[x] = true; slk[x] = INF;
if (fl[x] != -1) return (vr[q[r++] = fl[x]] = true);
   if (d[i][s] < inf && d[i + 1][t] > d[i][s] + c) {
   d[i + 1][t] = d[i][s] + c;
                                                               while (x != -1) swap(x, fr[fl[x] = pre[x]]);
    prv[i + 1][t] = s; prve[i + 1][t] = j;
                                                               return false;
  }
 }
                                                              while (true) {
llf mmc = inf; int st = -1;
                                                               while (l < r)</pre>
                                                                for (int x = 0, y = q[l++]; x < n; ++x) if (!vl[x])
 for (int i = 0; i < n; i++) {</pre>
                                                                 if (chmin(slk[x], hl[x] + hr[y] - w[x][y]))
 llf avg = -inf;
  for (int k = 0; k < n; k++) {
                                                                  if (pre[x] = y, !check(x)) return;
  if (d[n][i] < inf - eps)
                                                               lld d = ranges::min(slk);
                                                               for (int x = 0; x < n; ++x)
    avg = max(avg, (d[n][i] - d[k][i]) / (n - k));
   else avg = inf;
                                                                vl[x] ? hl[x] += d : slk[x] -= d;
                                                               for (int x = 0; x < n; ++x) if (vr[x]) hr[x] -= d;
                                                               for (int x = 0; x < n; ++x) if (!check(x)) return;
  if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
 if (st == -1) return inf;
                                                             KM(int n_, const auto &w) : n(n_), ans(0),
vector<int> vst(n), eid, cycle, rho;
                                                              hl(n), hr(n), fl(n, -1), fr(fl), pre(n), q(n) {
for (int i = n; !vst[st]; st = prv[i--][st]) {
                                                              for (int i = 0; i < n; ++i) hl[i]=ranges::max(w[i]);</pre>
 vst[st]++; eid.emplace_back(prve[i][st]);
  rho.emplace_back(st);
                                                              for (int i = 0; i < n; ++i) bfs(w, i);</pre>
                                                              for (int i = 0; i < n; ++i) ans += w[i][fl[i]];</pre>
```

while (not bfs.empty() and not lv[ed]) { int u = bfs.front(); bfs.pop();

bfs.push(e.to), lv[e.to] = lv[u] + 1;

for (auto e: G[u]) if (e.cap >> k and !lv[e.to])

```
return lv[ed];
}; // find maximum perfect matching
                                                                                Cap DFS(int u, Cap f = numeric_limits<Cap>::max()) {
// To obtain the max match of exactly K edges for
// K = 1 ... N, initialize hl[i] = INF and bfs from all
                                                                                 if (u == ed) return f;
// unmatched right part point (fr[i] == -1)
                                                                                 Cap ret = 0;
        Flow Models
                                                                                 for (auto &i = idx[u]; i < G[u].size(); ++i) {</pre>
                                                                                   auto &[to, rev, cap] = G[u][i];
· Maximum/Minimum flow with lower bound / Circulation problem
                                                                                   if (cap <= 0 or lv[to] != lv[u] + 1) continue;</pre>
  1. Construct super source {\cal S} and sink {\cal T}.
                                                                                   Cap nf = DFS(to, min(f, cap));
     For each edge (x, y, l, u), connect x \to y with capacity u - l.
                                                                                   ret += nf; cap -= nf; f -= nf;
     For each vertex v, denote by in(v) the difference between the sum of
     incoming lower bounds and the sum of outgoing lower bounds.
                                                                                   G[to][rev].cap += nf;
  4. If in(v) > 0, connect S \rightarrow v with capacity in(v), otherwise, connect
                                                                                   if (f == 0) return ret;
     v 	o T with capacity -in(v).
     - To maximize, connect t \to s with capacity \infty (skip this in circu-
                                                                                 if (ret == 0) lv[u] = 0;
       lation problem), and let f be the maximum flow from S to T. If
                                                                                 return ret;
       f 
eq \sum_{v \in V, in(v) > 0} in(v), there's no solution. Otherwise, the maxi-
       mum flow from s to t is the answer. Also, f is a mincost valid flow.
      To minimize, let f be the maximum flow from S to T. Connect t \to s with capacity \infty and let the flow from S to T be f'. If f+f' \ne \sum_{v \in V, in(v)>0} in(v), there's no solution. Otherwise, f' is the answer.
                                                                               public:
                                                                                void init(int n_) { G.assign(n = n_, vector<E>()); }
                                                                                void add_edge(int u, int v, Cap c) {
  5. The solution of each edge e is l_e+f_e, where f_e corresponds to the flow
                                                                                 G[u].push_back({v, int(G[v].size()), c});
G[v].push_back({u, int(G[u].size())-1, 0});
     of edge e on the graph.

    Construct minimum vertex cover from maximum matching M on bipartite

  graph(X, Y)
                                                                                Cap max_flow(int st_, int ed_) {
  1. Redirect every edge: y \to x if (x,y) \in M, x \to y otherwise. 2. DFS from unmatched vertices in X.
                                                                                 st = st_, ed = ed_; Cap ret = 0;
  3. x \in X is chosen iff x is unvisited; y \in Y is chosen iff y is visited.
                                                                                 for (int i = 63; i >= 0; --i)
· Minimum cost cyclic flow
                                                                                   while (BFS(i)) ret += DFS(st);
  1. Consruct super source {\cal S} and sink {\cal T}
                                                                                 return ret;
  2. For each edge (x,y,c), connect x \to y with (cost,cap)=(c,1) if c>0, otherwise connect y \to x with (cost,cap)=(-c,1)
                                                                               }; // test @ luogu P3376
4.5 HLPP [198e4e]
     For each edge with c<0, sum these cost as K, then increase d(y) by 1,
     decrease d(x) by 1
  4. For each vertex v with d(v) > 0, connect S \rightarrow v with (cost, cap) =
                                                                               template <typename T> struct HLPP {
     (0, d(v))
                                                                                struct Edge { int to, rev; T flow, cap; };
  5. For each vertex v with d(v) < 0, connect v \rightarrow T with (cost, cap) =
                                                                                int n, mx; vector<vector<Edge>> adj; vector<T> excess;
                                                                                vector<int> d, cnt, active; vector<vector<int>> B;
  6. Flow from S to T, the answer is the cost of the flow C+K

    Maximum density induced subgraph

                                                                                void add_edge(int u, int v, int f) {
                                                                                 Edge a{v, (int)size(adj[v]), 0, f};
Edge b{u, (int)size(adj[u]), 0, 0};
  1. Binary search on answer, suppose we're checking answer {\cal T}
     Construct a max flow model, let K be the sum of all weights
  3. Connect source s \to v, v \in G with capacity K
                                                                                 adj[u].push_back(a), adj[v].push_back(b);
  4. For each edge (u,v,w) in G, connect u 	o v and v 	o u with capacity w
  5. For v \in \mathit{G}, connect it with sink v \rightarrow t with capacity K + 2T
                                                                                void enqueue(int v) {
     \left(\sum_{e \in E(v)} w(e)\right) - 2w(v)
                                                                                 if (!active[v] && excess[v] > 0 && d[v] < n) {</pre>
  6. \stackrel{\searrow}{T} is a valid answer if the maximum flow f < K|V|
                                                                                   mx = max(mx, d[v]);
 Minimum weight edge cover
                                                                                   B[d[v]].push_back(v); active[v] = 1;
  1. For each v \in V create a copy v', and connect u' \to v' with weight
                                                                                 }
     w(u,v).
  2. Connect v -
                 	o v' with weight 2\mu(v) , where \mu(v) is the cost of the cheap-
     est edge incident to v.
                                                                                void push(int v, Edge &e) {
  3. Find the minimum weight perfect matching on G^\prime
                                                                                 T df = min(excess[v], e.cap - e.flow);
• Project selection cheat sheet: S,T 分別代表 0,1 側,最小化總花費。
   i 為 0 時花費 c
                                                                                 if (df <= 0 || d[v] != d[e.to] + 1) return;</pre>
                                  (i, T, c)
                                                                                 e.flow += df, adj[e.to][e.rev].flow -= df;
   i 為 1 時花費 c
                                  (S, i, c)
   i \in I 有任何一個為 0 時花費 c i \in I 有任何一個為 1 時花費 c
                                                                                 excess[e.to] += df, excess[v] -= df;
                                  (i, w, \infty), (w, T, c)
                                   S, w, c), (w, i, \infty)
                                                                                 enqueue(e.to);
   i 為 0 時得到 c
                                  直接得到 c; (S, i, c)
   i 為 1 時得到 c
                                  直接得到 c; (i, T, c)
                                                                                void gap(int k) {
   i 為 0 , j 為 1 時花費 c
                                  (i, j, c)
   i, j 不同時花費 c
                                                                                 for (int v = 0; v < n; v++) if (d[v] >= k)
                                  (i, j, c), (j, i, c)
   i, j 同時是 0 時得到 c
                                  直接得到 c; (S, w, c), (w, i, \infty), (w, j, \infty)
                                                                                   cnt[d[v]]--, d[v] = n, cnt[d[v]]++;
    i,j 同時是 1 時得到 c
                                  直接得到 c; (i, w, \infty), (j, w, \infty), (w, T, c)

    Submodular functions minimization

                                                                                void relabel(int v) {
 cnt[d[v]]--; d[v] = n;
                                                                                 for (auto e : adj[v])
                                                                                   if (e.cap > e.flow) d[v] = min(d[v], d[e.to] + 1);
                                                                                 cnt[d[v]]++; enqueue(v);
  - If \theta_i(1) \geq \theta_i(0), add edge (S, i, \theta_i(1) - \theta_i(0)) and \theta_i(0) to answer; other-
    wise, (i, T, \theta_i(0) - \theta_i(1)) and \theta_i(1)
 where, \{i, j, v_i(0) - v_{ij}\} and \{i, j, \phi_{ij}(0, 1) + \phi_{ij}(1, 0) - \phi_{ij}(0, 0) - \phi_{ij}(1, 1)\}.

Denote x_{ijk} as helper nodes. Let P = \psi_{ijk}(0, 0, 0) + \psi_{ijk}(0, 1, 1) + \psi_{ijk}(1, 0, 1) + \psi_{ijk}(1, 1, 0) - \psi_{ijk}(0, 0, 1) - \psi_{ijk}(0, 1, 0) - \psi_{ijk}(1, 1, 1). Add P to answer. If P \geq 0, add edges \{i, x_{ijk}, P\},
                                                                                void discharge(int v) {
                                                                                 for (auto &e : adj[v])
                                                                                   if (excess[v] > 0) push(v, e);
    (x_{ijk}, P), (x_{ijk}, P), (x_{ijk}, P), (x_{ijk}, T, P); otherwise (x_{ijk}, i, -P), (x_{ijk}, j, -P), (x_{ijk}, k, -P), (x_{ijk}, k, -P).
                                                                                   else break;
                                                                                 if (excess[v] <= 0) return;</pre>
    The minimum cut of this graph will be the the minimum value of the
                                                                                 if (cnt[d[v]] == 1) gap(d[v]);
    function above.
                                                                                 else relabel(v);
4.4 Dinic [32c53e]
                                                                                T max_flow(int s, int t) {
template <typename Cap = int64_t> class Dinic {
                                                                                 for (auto &e : adj[s]) excess[s] += e.cap;
private:
 struct E { int to, rev; Cap cap; }; int n, st, ed;
                                                                                 cnt[0] = n; enqueue(s); active[t] = 1;
 vector<vector<E>> G; vector<size_t> lv, idx;
                                                                                 for (mx = 0; mx >= 0;)
                                                                                   if (!B[mx].empty())
 bool BFS(int k) {
  lv.assign(n, 0); idx.assign(n, 0);
                                                                                    int v = B[mx].back(); B[mx].pop_back();
  queue<int> bfs; bfs.push(st); lv[st] = 1;
                                                                                    active[v] = 0; discharge(v);
```

} else --mx;

return excess[t];

HLPP(int \_n) : n(\_n), adj(n), excess(n), d(n), cnt(n + 1), active(n), B(n) {}

```
|};
                                                                return {c, w};
 4.6
      Global Min-Cut [ae7013]
                                                               }
 void add_edge(auto &w, int u, int v, int c) {
 w[u][v] += c; w[v][u] += c; }
                                                              4.9
                                                                    Dijkstra Cost Flow [d0cfd9]
 auto phase(const auto &w, int n, vector<int> id) {
                                                              template <typename F, typename C> class MCMF {
 vector<lld> g(n); int s = -1, t = -1;
                                                               static constexpr F INF_F = numeric_limits<F>::max();
 while (!id.empty()) {
                                                               static constexpr C INF_C = numeric_limits<C>::max();
   int c = -1;
                                                               struct E { int to, r; F f; C c; };
   for (int i : id) if (c == -1 || g[i] > g[c]) c = i;
                                                               vector<vector<E>> g; vector<pair<int, int>> f;
   s = t; t = c;
                                                               vector<F> up; vector<C> d, h;
   id.erase(ranges::find(id, c));
                                                               optional<pair<F, C>> step(int S, int T) {
   for (int i : id) g[i] += w[c][i];
                                                                priority_queue<pair<C, int>> q;
                                                                q.emplace(d[S] = 0, S), up[S] = INF_F;
 return tuple{s, t, g[t]};
                                                                while (not q.empty()) {
                                                                 auto [l, u] = q.top(); q.pop();
 lld mincut(auto w, int n) {
                                                                  if (up[u] == 0 or l != -d[u]) continue;
 lld cut = numeric_limits<lld>::max();
                                                                 for (int i = 0; i < int(g[u].size()); ++i) {</pre>
 vector<int> id(n); iota(all(id), 0);
for (int i = 0; i < n - 1; ++i) {</pre>
                                                                  auto e = g[u][i]; int v = e.to;
                                                                  auto nd = d[u] + e.c + h[u] - h[v];
  auto [s, t, gt] = phase(w, n, id);
                                                                  if (e.f <= 0 or d[v] <= nd) continue;</pre>
  id.erase(ranges::find(id, t));
                                                                  f[v] = {u, i}; up[v] = min(up[u], e.f);
   cut = min(cut, gt);
                                                                  q.emplace(-(d[v] = nd), v);
  for (int j = 0; j < n; ++j)</pre>
   w[s][j] += w[t][j], w[j][s] += w[j][t];
                                                                if (d[T] == INF_C) return nullopt;
 return cut;
                                                                for (size_t i = 0; i < d.size(); ++i) h[i] += d[i];</pre>
 \} // O(V^3), can be O(VE + V^2 \log V)?
                                                                for (int i = T; i != S; i = f[i].first) {
       GomoryHu Tree [245ce3]
                                                                 auto &eg = g[f[i].first][f[i].second];
 auto GomoryHu(int n, const auto &flow) {
                                                                 eg.f -= up[T]; g[eg.to][eg.r].f += up[T];
 vector<tuple<int, int, int>> rt; vector<int> g(n);
for (int i = 1; i < n; ++i) {</pre>
                                                                return pair{up[T], h[T]};
  int t = g[i]; auto f = flow;
  rt.emplace_back(f.max_flow(i, t), i, t);
                                                              public:
  f.walk(i); // bfs from i use edges with .cap > 0
                                                               MCMF(int n) : g(n), f(n), up(n), d(n, INF_C) {}
  for (int j = i + 1; j < n; ++j)</pre>
                                                               void add_edge(int s, int t, F c, C w) {
   if (g[j]==t && f.connect(j)) g[j] = i;
                                                                g[s].emplace_back(t, int(g[t].size()), c, w);
                                                                g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
 return rt;
} // for our dinic:
// void walk(int) { BFS(0); }
                                                               pair<F, C> solve(int a, int b) {
                                                                h.assign(g.size(), 0);
 // bool connect(int i) { return lv[i]; }
                                                                F c = 0; C w = 0;
      MCMF [0df510]
                                                                while (auto r = step(a, b)) {
template <typename F, typename C> class MCMF {
    static constexpr F INF_F = numeric_limits<F>::max();
                                                                 c += r->first, w += r->first * r->second;
                                                                 fill(d.begin(), d.end(), INF_C);
 static constexpr C INF_C = numeric_limits<C>::max();
 struct E { int to, r; F f; C c; };
                                                                return {c, w};
                                                               7
 vector<vector<E>> g; vector<pair<int, int>> f;
 vector<int> inq; vector<F> up; vector<C> d;
                                                              4.10 Min Cost Circulation [ea0477]
 optional<pair<F, C>> step(int S, int T) {
   queue<int> q;
                                                              template <typename F, typename C>
   for (q.push(S), d[S] = 0, up[S] = INF_F;
  not q.empty(); q.pop()) {
                                                              struct MinCostCirculation {
                                                               struct ep { int to; F flow; C cost; };
    int u = q.front(); inq[u] = false;
                                                               int n; vector<int> vis; int visc;
    if (up[u] == 0) continue;
for (int i = 0; i < int(g[u].size()); ++i) {</pre>
                                                               vector<int> fa, fae; vector<vector<int>> g;
                                                               vector<ep> e; vector<C> pi;
     auto e = g[u][i]; int v = e.to;
                                                               MinCostCirculation(int n_) : n(n_), vis(n), visc(0), g
     if (e.f <= 0 or d[v] <= d[u] + e.c) continue;</pre>
                                                                   (n), pi(n) {}
     d[v] = d[u] + e.c; f[v] = {u, i};
                                                               void add_edge(int u, int v, F fl, C cs) {
     up[v] = min(up[u], e.f);
                                                                g[u].emplace_back((int)e.size());
     if (not inq[v]) q.push(v);
                                                                e.emplace_back(v, fl, cs);
     inq[v] = true;
                                                                g[v].emplace_back((int)e.size());
                                                                e.emplace_back(u, 0, -cs);
   }
   if (d[T] == INF_C) return nullopt;
                                                               C phi(int x) {
   for (int i = T; i != S; i = f[i].first) {
                                                                if (fa[x] == -1) return 0;
                                                                if (vis[x] == visc) return pi[x];
    auto &eg = g[f[i].first][f[i].second];
    eg.f -= up[T]; g[eg.to][eg.r].f += up[T];
                                                                vis[x] = visc;
                                                                return pi[x] = phi(fa[x]) - e[fae[x]].cost;
  return pair{up[T], d[T]};
                                                               int lca(int u, int v) {
public:
                                                                for (; u != -1 || v != -1; swap(u, v)) if (u != -1) {
 MCMF(int n) : g(n),f(n),inq(n),up(n),d(n,INF_C) {}
                                                                 if (vis[u] == visc) return u;
 void add_edge(int s, int t, F c, C w) {
                                                                 vis[u] = visc; u = fa[u];
   g[s].emplace_back(t, int(g[t].size()), c, w);
  g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
                                                                return -1;
 pair<F, C> solve(int a, int b) {
                                                               void pushflow(int x, C &cost) {
  F c = 0; C w = 0;
                                                                int v = e[x ^1].to, u = e[x].to; ++visc;
  while (auto r = step(a, b)) {
                                                                if (int w = lca(u, v); w == −1) {
    c += r->first, w += r->first * r->second;
                                                                 while (v != -1)
    ranges::fill(inq, false); ranges::fill(d, INF_C);
                                                                  swap(x ^= 1, fae[v]), swap(u, fa[v]), swap(u, v);
```

```
} else {
   int z = u, dir = 0; F f = e[x].flow;
   vector<int> cyc = {x};
   for (int d : {0, 1})
    for (int i = (d ? u : v); i != w; i = fa[i]) {
     cyc.push_back(fae[i] ^ d);
     if (chmin(f, e[fae[i] ^ d].flow)) z = i, dir = d;
   for (int i : cyc) {
    e[i].flow -= f; e[i ^ 1].flow += f;
    cost += f * e[i].cost;
   if (dir) x ^= 1, swap(u, v);
   while (u != z)
    swap(x ^= 1, fae[v]), swap(u, fa[v]), swap(u, v);
 void dfs(int u) {
  vis[u] = visc;
  for (int i : g[u])
   if (int v = e[i].to; vis[v] != visc and e[i].flow)
    fa[v] = u, fae[v] = i, dfs(v);
 C simplex() {
 fa.assign(g.size(), -1); fae.assign(g.size(), -1);
C cost = 0; ++visc; dfs(0);
for (int fail = 0; fail < ssize(e); )</pre>
   for (int i = 0; i < ssize(e); i++)</pre>
    if (e[i].flow and e[i].cost < phi(e[i ^ 1].to) -</pre>
    phi(e[i].to))
     fail = 0, pushflow(i, cost), ++visc;
    else ++fail;
  return cost;
4.11 General Matching [5f2293]
struct Matching {
 queue<int> q; int ans, n;
 vector<int> fa, s, v, pre, match;
 int Find(int u) {
  return u == fa[u] ? u : fa[u] = Find(fa[u]); }
 int LCA(int x, int y) {
  static int tk = 0; tk++; x = Find(x); y = Find(y);
  for (;; swap(x, y)) if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
  }
 void Blossom(int x, int y, int l) {
  for (; Find(x) != l; x = pre[y]) {
   pre[x] = y, y = match[x];
   if (s[y] == 1) q.push(y), s[y] = 0;
   for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
  }
 bool Bfs(auto &&g, int r) {
  iota(all(fa), 0); ranges::fill(s, -1);
  q = queue<int>(); q.push(r); s[r] = 0;
  for (; !q.empty(); q.pop()) {
   for (int x = q.front(); int u : g[x])
    if (s[u] == -1) {
     if (pre[u] = x, s[u] = 1, match[u] == n) {
      for (int a = u, b = x, last;
        b != n; a = last, b = pre[a])
       last = match[b], match[b] = a, match[a] = b;
      return true;
     q.push(match[u]); s[match[u]] = 0;
    } else if (!s[u] && Find(u) != Find(x)) {
     int l = LCA(u, x);
     Blossom(x, u, l); Blossom(u, x, l);
    }
  }
  return false;
 Matching(auto &&g) : ans(0), n(int(g.size())),
 fa(n+1), s(n+1), v(n+1), pre(n+1, n), match(n+1, n) {
  for (int x = 0; x < n; ++x)</pre>
   if (match[x] == n) ans += Bfs(g, x);
 } // match[x] == n means not matched
}; // test @ yosupo judge
```

```
Weighted Matching [900530]
4.12
#define pb emplace_back
#define REP(i, l, r) for (int i=(l); i<=(r); ++i)
struct WeightGraph { // 1-based
 static const int inf = INT_MAX;
 struct edge { int u, v, w; }; int n, nx;
 vector<int> lab; vector<vector<edge>> g;
 vector<int> slack, match, st, pa, S, vis;
 vector<vector<int>> flo, flo_from; queue<int> q;
WeightGraph(int n_-): n(n_-), nx(n * 2), lab(nx + 1), g(nx + 1, vector < edge > (nx + 1), slack(nx + 1),
  flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
  match = st = pa = S = vis = slack;
 REP(u, 1, n) REP(v, 1, n) g[u][v] = \{u, v, 0\};
 int ED(edge e) {
 return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
 void update_slack(int u, int x, int &s) {
  if (!s || ED(g[u][x]) < ED(g[s][x])) s = u; }</pre>
 void set_slack(int x) {
  slack[x] = 0;
  REP(u, 1, n)
   if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
    update_slack(u, x, slack[x]);
 void q_push(int x) {
 if (x \le n) q.push(x);
  else for (int y : flo[x]) q_push(y);
 void set_st(int x, int b) {
  st[x] = b;
  if (x > n) for (int y : flo[x]) set_st(y, b);
 vector<int> split_flo(auto &f, int xr) {
 auto it = find(all(f), xr);
  if (auto pr = it - f.begin(); pr % 2 == 1)
  reverse(1 + all(f)), it = f.end() - pr;
  auto res = vector(f.begin(), it);
  return f.erase(f.begin(), it), res;
 void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  int xr = flo_from[u][g[u][v].u];
  auto &f = flo[u], z = split_flo(f, xr);
  REP(i, 0, int(z.size())-1) set_match(z[i], z[i ^ 1]);
  set_match(xr, v); f.insert(f.end(), all(z));
 void augment(int u, int v) {
  for (;;) {
   int xnv = st[match[u]]; set_match(u, v);
   if (!xnv) return;
   set_match(v = xnv, u = st[pa[xnv]]);
 int lca(int u, int v) {
  static int t = 0; ++t;
  for (++t; u || v; swap(u, v)) if (u) {
  if (vis[u] == t) return u;
   vis[u] = t; u = st[match[u]];
   if (u) u = st[pa[u]];
  return 0;
 void add_blossom(int u, int o, int v) {
  int b = int(find(n + 1 + all(st), 0) - begin(st));
 lab[b] = 0, S[b] = 0; match[b] = match[o];
vector<int> f = {0};
  for (int x : {u, v}) {
   for (int y; x != o; x = st[pa[y]])
    f.pb(x), f.pb(y = st[match[x]]), q_push(y);
   reverse(1 + all(f));
  flo[b] = f; set_st(b, b);
  REP(x, 1, nx) g[b][x].w = g[x][b].w = 0;
  REP(x, 1, n) flo_from[b][x] = 0;
  for (int xs : flo[b]) {
   REP(x, 1, nx)
    if (g[b][x].w == 0 \mid \mid ED(g[xs][x]) < ED(g[b][x]))
     g[b][x] = g[xs][x], g[x][b] = g[x][xs];
   REP(x, 1, n)
    if (flo_from[xs][x]) flo_from[b][x] = xs;
```

```
void set_edge(int u, int v, int w) {
 set_slack(b);
                                                                          g[u][v].w = g[v][u].w = w; }
}
                                                                       5
5.1
                                                                              Math
void expand_blossom(int b) {
 for (int x : flo[b]) set_st(x, x);
                                                                              Common Bounds
 int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
                                                                         n | 2 3 4 5 6 7 8 9 20 50 100 n | 100 le3 le6 le9 le12 le15
 for (int x : split_flo(flo[b], xr)) {
                                                                        \overline{p(n)} 2 3 5 7 11 15 22 30 627 2e5 2e8 \overline{d(i)} 12 32 240 1344 6720 26880 103680
  if (xs == -1) { xs = x; continue; }
                                                                                                         9
                                                                                                                10 11 12 13 14
                                                                                                   8
  pa[xs] = g[x][xs].u; S[xs] = 1, S[x] = 0;
                                                                        \binom{2n}{n} 2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 le7 4e7 1.5e8
  slack[xs] = 0; set_slack(x); q_push(x); xs = -1;
                                                                         n 2345 6 7
                                                                                            8
                                                                                                 9
                                                                                                       10
                                                                                                            11 12 13
                                                                        \frac{1}{B_n} 2 5 15 52 203 877 4140 21147 115975 7e5 4e6 3e7
 for (int x : flo[b])
                                                                        5.2 Equations
  if (x == xr) S[x] = 1, pa[x] = pa[b];
                                                                        Stirling Number of the First Kind
  else S[x] = -1, set_slack(x);
                                                                        S_1(n,k) counts the number of permutations of n elements with k disjoint
 st[b] = 0;
                                                                        cycles.
                                                                        • S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)
bool on_found_edge(const edge &e) {
                                                                        • S_1(n,i) = [x^i] \left(\prod_{i=0}^{n-1} (x+i)\right), use D&Q and taylor shift.
 if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
                                                                        • S_1(i,k) = \frac{i!}{k!} \left[ x^i \right] \left( \sum_{j \ge 1} \frac{x^j}{j} \right)
  int nu = st[match[v]]; pa[v] = e.u; S[v] = 1;
  slack[v] = slack[nu] = 0; S[nu] = 0; q_push(nu);
                                                                        Stirling Number of the Second Kind
 } else if (S[v] == 0) {
                                                                        S_2(n,k) counts the number of ways to partition a set of n elements into k
  if (int o = lca(u, v)) add_blossom(u, o, v);
                                                                        nonempty sets.
                                                                        • S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)
  else return augment(u, v), augment(v, u), true;
                                                                        • S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}
                                                                        \cdot \ S_2(i,k) = rac{i!}{k!} [x^i] \, (e^x - 1)^k Derivatives/Integrals
 return false;
                                                                        bool matching() {
 ranges::fill(S, -1); ranges::fill(slack, 0);
                                                                        \frac{d}{dx}\sin^{-1}x = \frac{1}{\sqrt{1-x^2}} \begin{vmatrix} \frac{d}{dx}\cos^{-1}x = -\frac{1}{\sqrt{1-x^2}} \\ \frac{d}{dx}\tan x = 1 + \tan^2 x \end{vmatrix} \int \tan ax = -\frac{\ln|\cos ax|}{a} \begin{vmatrix} \frac{d}{dx}\tan^{-1}x = \frac{1}{1+x^2} \\ \frac{d}{dx}\tan x = \frac{1}{1+x^2} \end{vmatrix}
 q = queue<int>();
 REP(x, 1, nx) if (st[x] == x \&\& !match[x])
  pa[x] = 0, S[x] = 0, q_push(x);
                                                                         \int_{0}^{ax} e^{-x^{2}} = \frac{\sqrt{\pi}}{2} \operatorname{erf}(x) \left| \int_{0}^{ax} x \, dx = \frac{e^{ax}}{a^{2}} (ax - 1) \right|
 if (q.empty()) return false;
 for (;;) {
                                                                           \int \sqrt{a^2 + x^2} = \frac{1}{2} \left( x \sqrt{a^2 + x^2} + a^2 \operatorname{asinh}(x/a) \right)
  while (q.size()) {
   int u = q.front(); q.pop();
                                                                        Extended Euler
                                                                       a^b \equiv \begin{cases} a^{(b \mod \varphi(m)) + \varphi(m)} & \text{if } (a,m) \neq 1 \wedge b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m
    if (S[st[u]] == 1) continue;
   REP(v, 1, n)
     if (g[u][v].w > 0 && st[u] != st[v]) {
                                                                        Pentagonal Number Theorem
      if (ED(g[u][v]) != 0)
                                                                        \prod_{n=1}^{\infty} (1 - x^n) = \sum_{k=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2} = (\sum p(n)x^n)^{-1}
       update_slack(u, st[v], slack[st[v]]);
                                                                        5.3 Integer Division* [cd017d]
      else if (on_found_edge(g[u][v])) return true;
                                                                       lld fdiv(lld a, lld b)
{ return a / b - (a % b && (a < 0) ^ (b < 0)); }</pre>
     }
                                                                        lld cdiv(lld a, lld b)
  int d = inf;
  REP(b, n + 1, nx) if (st[b] == b && S[b] == 1)
d = min(d, lab[b] / 2);
                                                                         return a / b + (a % b && (a < 0) ^ (b > 0)); }
                                                                        5.4 FloorSum [fb5917]
  REP(x, 1, nx)
                                                                       // @param n `n < 2^32
                                                                        // @param m `1 <= m < 2^32
   if (int s = slack[x]; st[x] == x && s && S[x] <= 0)</pre>
                                                                        // @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
     d = min(d, ED(g[s][x]) / (S[x] + 2));
  REP(u, 1, n)
                                                                        llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
   if (S[st[u]] == 1) lab[u] += d;
                                                                         llu ans = 0;
   else if (S[st[u]] == 0) {
                                                                         while (true) {
                                                                          if (a >= m) ans += n*(n-1)/2 * (a/m), a %= m;
     if (lab[u] <= d) return false;</pre>
     lab[u] -= d;
                                                                          if (b >= m) ans += n * (b/m), b %= m;
                                                                          if (llu y_max = a * n + b; y_max >= m) {
  n = (llu)(y_max / m), b = (llu)(y_max % m);
  REP(b, n + 1, nx) if (st[b] == b && S[b] >= 0)
   lab[b] += d * (2 - 4 * S[b]);
                                                                            swap(m, a);
  REP(x, 1, nx)
                                                                          } else break;
   if (int s = slack[x]; st[x] == x &&
      s \&\& st[s] != x \&\& ED(g[s][x]) == 0)
                                                                         return ans;
     if (on_found_edge(g[s][x])) return true;
  REP(b, n + 1, nx)
                                                                        lld floor_sum(lld n, lld m, lld a, lld b) {
   if (st[b] == b && S[b] == 1 && lab[b] == 0)
                                                                         llu ans = 0;
     expand_blossom(b);
                                                                         if (a < 0) {
                                                                          llu a2 = (a \% m + m), d = (a2 - a) / m;
 return false;
                                                                          ans -= 1ULL * n * (n - 1) / 2 * d; a = a2;
}
pair<lld, int> solve() {
 ranges::fill(match, 0);
                                                                          llu b2 = (b \% m + m), d = (b2 - b) / m;
 REP(u, 0, n) st[u] = u, flo[u].clear();
                                                                          ans -= 1ULL * n * d; b = b2;
 int w_max = 0;
 REP(u, 1, n) REP(v, 1, n) {
                                                                         return ans + floor_sum_unsigned(n, m, a, b);
  flo_from[u][v] = (u == v ? u : 0);
                                                                        5.5 ModMin [2c021c]
  w_max = max(w_max, g[u][v].w);
                                                                        // min{k | l <= ((ak) mod m) <= r}
 REP(u, 1, n) lab[u] = w_max;
                                                                       optional<llu> mod_min(u32 a, u32 m, u32 l, u32 r) {
 int n_matches = 0; lld tot_weight = 0;
while (matching()) ++n_matches;
                                                                         if (a == 0) return l ? nullopt : optional{0};
                                                                         if (auto k = llu(l + a - 1) / a; k * a <= r)
 REP(u, 1, n) if (match[u] \&\& match[u] < u)
                                                                          return k;
  tot_weight += g[u][match[u]].w;
                                                                         auto b = m / a, c = m % a;
 return make_pair(tot_weight, n_matches);
                                                                         if (auto y = mod_min(c, a, a - r % a, a - l % a))
                                                                          return (l + *y * c + a - 1) / a + *y * b;
```

```
return nullopt;
                                                                if (q \& 1) r = M(r, e);
                                                               return int(r.first); // sqrt(n) mod P where P is prime
5.6 Floor Monoid Product [416e89]
                                                              5.11 FWT [88a937]
/* template <typename T>
T brute(llu a, llu b, llu c, llu n, T U, T R) {
                                                             /* or convolution:
                                                              * x = (x0, x0+x1), inv = (x0, x1-x0)  w/o final div
 for (llu\ i = 1, l = 0; i <= n; i++, res = res * R)
                                                               * and convolution:
                                                              * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
void fwt(int x[], int N, bool inv = false) {
 for (llu \ r = (a*i+b)/c; \ l < r; ++l) \ res = res * U;
 return res;
                                                               for (int d = 1; d < N; d <<= 1)</pre>
} */
template <typename T>
T euclid(llu a, llu b, llu c, llu n, T U, T R) {
                                                                for (int s = 0; s < N; s += d * 2)
                                                                 for (int i = s; i < s + d; i++) {</pre>
                                                                  int j = i + d, ta = x[i], tb = x[j];
 if (!n) return T{};
 if (b >= c)
                                                                  x[i] = add(ta, tb); x[j] = sub(ta, tb);
  return mpow(U, b / c) * euclid(a, b % c, c, n, U, R);
                                                               if (!inv) return;
 if (a >= c)
  return euclid(a % c, b, c, n, U, mpow(U, a / c) * R);
                                                               const int invn = modinv(N);
 llu m = (u128(a) * n + b) / c;
                                                               for (int i = 0; i < N; i++) x[i] = mul(x[i], invn);</pre>
 if (!m) return mpow(R, n);
 return mpow(R, (c - b - 1) / a) * U
                                                              5.12 Packed FFT [0a6af5]
  * euclid(c, (c - b - 1) % a, a, m - 1, R, U)
                                                              VL convolution(const VI &a, const VI &b) {
  * mpow(R, n - (u128(c) * m - b - 1) / a);
                                                               if (a.empty() || b.empty()) return {};
                                                               const int sz = bit_ceil(a.size() + b.size() - 1);
// time complexity is O(log max(a, b, c))
                                                               // Should be able to handle N <= 10^5, C <= 10^4
// UUUU R UUUUU R ... UUU R 共 N 個 R,最後一個必是 R
                                                               vector<P> v(sz);
// 一直到第 k 個 R 前總共有 (ak+b)/c 個 U
                                                               for (size_t i = 0; i < a.size(); ++i) v[i].RE(a[i]);</pre>
5.7 ax+by=gcd [d0cbdd]
                                                               for (size_t i = 0; i < b.size(); ++i) v[i].IM(b[i]);</pre>
// ax+ny = 1, ax+ny == ax == 1 \ (mod \ n)
                                                               fft(v.data(), sz, /*inv=*/false);
void exgcd(lld x, lld y, lld &g, lld &a, lld &b) {
                                                               auto rev = v; reverse(1 + all(rev));
 if (y == 0) g = x, a = 1, b = 0;
                                                               for (int i = 0; i < sz; ++i) {
 else exgcd(y, x \% y, g, b, a), b -= (x / y) * a;
                                                                P A = (v[i] + conj(rev[i])) / P(2, 0);
                                                               P B = (v[i] - conj(rev[i])) / P(0, 2);
5.8 Chinese Remainder [d69e74]
                                                               v[i] = A * B;
// please ensure r_i\in[0,m_i)
                                                               VL c(sz); fft(v.data(), sz, /*inv=*/true);
for (int i = 0; i < sz; ++i) c[i] = roundl(RE(v[i]));</pre>
bool crt(lld &m1, lld &r1, lld m2, lld r2) {
  if (m2 > m1) swap(m1, m2), swap(r1, r2);
                                                               return c;
  lld g, a, b; exgcd(m1, m2, g, a, b);
  if ((r2 - r1) % g != 0) return false;
  m2 /= g; lld D = (r2 - r1) / g % m2 * a % m2;
                                                             VI convolution_mod(const VI &a, const VI &b) {
 r1 += (D < 0 ? D + m2 : D) * m1; m1 *= m2;
                                                               if (a.empty() || b.empty()) return {};
                                                               const int sz = bit_ceil(a.size() + b.size() - 1);
  assert (r1 >= 0 && r1 < m1);
  return true;
                                                               vector<P> fa(sz), fb(sz);
                                                               for (size_t i = 0; i < a.size(); ++i)</pre>
                                                                fa[i] = P(a[i] & ((1 << 15) - 1), a[i] >> 15);
5.9 DiscreteLog [86e463]
                                                               for (size_t i = 0; i < b.size(); ++i)
fb[i] = P(b[i] & ((1 << 15) - 1), b[i] >> 15);
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
                                                               fft(fa.data(), sz); fft(fb.data(), sz);
 // x^? \setminus equiv y \pmod{M}
 Int t = 1, c = 0, g = 1;
                                                               auto rfa = fa; reverse(1 + all(rfa));
                                                               for (int i = 0; i < sz; ++i) fa[i] *= fb[i];</pre>
 for (Int M_{-} = M; M_{-} > 0; M_{-} >>= 1) g = g * x % M;
                                                               for (int i = 0; i < sz; ++i) fb[i] *= conj(rfa[i]);</pre>
 for (g = gcd(g, M); t % g != 0; ++c) {
                                                               fft(fa.data(), sz, true); fft(fb.data(), sz, true);
 if (t == y) return c;
                                                               vector<int> res(sz);
 t = t * x % M;
                                                               for (int i = 0; i < sz; ++i) {</pre>
                                                                lld A = (lld)roundl(RE((fa[i] + fb[i]) / P(2, 0)));
 if (y % g != 0) return -1;
                                                                lld C = (lld) roundl(IM((fa[i] - fb[i]) / P(0, 2)));
 t /= g, y /= g, M /= g;
                                                                lld B = (lld)roundl(IM(fa[i])); B %= p; C %= p;
 Int h = 0, gs = 1;

for (; h * h < M; ++h) gs = gs * x % M;
                                                                res[i] = (A + (B << 15) + (C << 30)) % p;
 unordered_map<Int, Int> bs;
                                                               return res;
 for (Int s = 0; s < h; bs[y] = ++s) y = y * x % M;
                                                              } // test @ yosupo judge with long double
 for (Int s = 0; s < M; s += h) {
                                                              5.13
                                                                    CRT for arbitrary mod [e4dde7]
 t = t * gs % M;
  if (bs.count(t)) return c + s + h - bs[t];
                                                              const int mod = 1000000007;
                                                              const int M1 = 985661441; // G = 3 for M1, M2, M3
                                                              const int M2 = 998244353;
 return -1;
                                                              const int M3 = 1004535809;
                                                              int superBigCRT(lld A, lld B, lld C) {
5.10 Quadratic Residue [f0baec]
                                                               static_assert (M1 < M2 && M2 < M3);</pre>
int get_root(int n, int P) { // ensure 0 <= n < p
if (P == 2 or n == 0) return n;</pre>
                                                               constexpr lld r12 = modpow(M1, M2-2, M2);
                                                               constexpr lld r13 = modpow(M1, M3-2, M3);
 auto check = [&](lld x) {
                                                               constexpr lld r23 = modpow(M2, M3-2, M3);
  return modpow(int(x), (P - 1) / 2, P); };
                                                               constexpr lld M1M2 = 1LL * M1 * M2 % mod;
 if (check(n) != 1) return -1;
                                                               B = (B - A + M2) * r12 % M2;
 mt19937 \text{ rnd}(7122); \text{ lld } z = 1, w;
                                                               C = (C - A + M3) * r13 % M3;
 while (check(w = (z * z - n + P) % P) != P - 1)
                                                               C = (C - B + M3) * r23 % M3;
 z = rnd() \% P;
                                                               return (A + B * M1 + C * M1M2) % mod;
 const auto M = [P, w](auto &u, auto &v) {
 auto [a, b] = u; auto [c, d] = v;
return make_pair((a * c + b * d % P * w) % P,
                                                              5.14 NTT / FFT [2ac7d2]
    (a * d + b * c) \% P);
                                                              template <int mod, int G, int maxn> struct NTT {
                                                               static_assert(maxn == (maxn & -maxn));
 pair<lld, lld> r(1, 0), e(z, 1);
                                                               int roots[maxn];
 for (int q = (P + 1) / 2; q; q >>= 1, e = M(e, e))
                                                              NTT () {
```

```
int r = modpow(G, (mod - 1) / maxn);
                                                              S Sqrt(const S &v) { // need: QuadraticResidue
                                                               assert(!v.empty() && v[0] != 0);
  for (int i = maxn >> 1; i; i >>= 1) {
  roots[i] = 1;
                                                               const int r = get_root(v[0]); assert(r != -1);
                                                               return Newton(v, r,
[](S &X, S &A, int sz) {
   for (int j = 1; j < i; j++)</pre>
   roots[i + j] = mul(roots[i + j - 1], r);
   r = mul(r, r);
                                                                 auto Y = X; Y.resize(sz / 2);
   // for (int j = 0; j < i; j++) // FFT (tested)
                                                                 auto B = Mul(A, Inv(Y), sz);
                                                                 for (int i = 0, inv2 = mod / 2 + 1; i < sz; i++)</pre>
     roots[i+j] = polar < llf > (1, PI * j / i);
                                                                  X[i] = mul(inv2, add(X[i], B[i])); });
 // n must be 2^k, and 0 \le F[i] \le mod
                                                              S Mul(auto &&a, auto &&b) {
const auto n = a.size() + b.size() - 1;
                                                               auto R = Mul(a, b, bit_ceil(n));
                                                               return R.resize(n), R;
   for (int k = n>>1; (j^=k) < k; k>>=1);
                                                              S MulT(S a, S b, size_t k) {
  for (int s = 1; s < n; s *= 2)
                                                               assert(b.size()); reverse(all(b)); auto R = Mul(a, b);
   for (int i = 0; i < n; i += s * 2)
                                                               R = vector(R.begin() + b.size() - 1, R.end());
    for (int j = 0; j < s; j++) {
                                                               return R.resize(k), R;
     int a = F[i+j], b = mul(F[i+j+s], roots[s+j]);
F[i+j] = add(a, b); F[i+j+s] = sub(a, b);
                                                              S Eval(const S &f, const S &x) {
                                                               if (f.empty()) return vector(x.size(), 0);
  if (!inv) return;
                                                               const int n = int(max(x.size(), f.size()));
  const int invn = modinv(n);
                                                               auto q = vector(n * 2, S(2, 1)); S ans(n);
  for (int i = 0; i < n; i++) F[i] = mul(F[i], invn);</pre>
                                                               fi(0, x.size()) q[i + n][1] = sub(0, x[i]);
  reverse(F + 1, F + n);
                                                               for (int i = n - 1; i > 0; i--)
                                                                q[i] = Mul(q[i << 1], q[i << 1 | 1]);
                                                               q[1] = MulT(f, Inv(q[1]), n);
                                                               for (int i = 1; i < n; i++) {
      Formal Power Series [c6b99a]
                                                                auto L = q[i << 1], R = q[i << 1 | 1];</pre>
#define fi(l, r) for (size_t i = (l); i < (r); i++)
                                                                q[i << 1 | 0] = MulT(q[i], R, L.size());</pre>
using S = vector<int>;
                                                                q[i << 1 | 1] = MulT(q[i], L, R.size());</pre>
auto Mul(auto a, auto b, size_t sz) {
a.resize(sz), b.resize(sz);
ntt(a.data(), sz); ntt(b.data(), sz);
                                                               for (int i = 0; i < n; i++) ans[i] = q[i + n][0];</pre>
fi(0, sz) a[i] = mul(a[i], b[i]);
                                                               return ans.resize(x.size()), ans;
return ntt(a.data(), sz, true), a;
                                                              pair<S, S> DivMod(const S &A, const S &B) {
S Newton(const S &v, int init, auto &&iter) {
                                                               assert(!B.empty() && B.back() != 0);
S Q = { init };
                                                               if (A.size() < B.size()) return {{}}, A};</pre>
                                                               const auto sz = A.size() - B.size() + 1;
for (int sz = 2; Q.size() < v.size(); sz *= 2) {</pre>
 S A{begin(v), begin(v) + min(sz, int(v.size()))};
                                                               S X = B; reverse(all(X)); X.resize(sz);
 A.resize(sz * 2), Q.resize(sz * 2);
                                                               S Y = A; reverse(all(Y)); Y.resize(sz);
                                                               S Q = Mul(Inv(X), Y);
 iter(Q, A, sz * 2); Q.resize(sz);
                                                               Q.resize(sz); reverse(all(Q)); X = Mul(Q, B); Y = A;
                                                               fi(0, Y.size()) Y[i] = sub(Y[i], X[i]);
return Q.resize(v.size()), Q;
                                                               while (Y.size() && Y.back() == 0) Y.pop_back();
                                                               while (Q.size() && Q.back() == 0) Q.pop_back();
S Inv(const S &v) { // v[0] != 0
return Newton(v, modinv(v[0]),
  [](S &X, S &A, int sz) {
                                                               return {Q, Y};
                                                               // empty means zero polynomial
                                                              int LinearRecursionKth(S a, S c, int64_t k) {
  ntt(X.data(), sz), ntt(A.data(), sz);
                                                               const auto d = a.size(); assert(c.size() == d + 1);
   for (int i = 0; i < sz; i++)</pre>
   X[i] = mul(X[i], sub(2, mul(X[i], A[i])));
                                                               const auto sz = bit_ceil(2 * d + 1), o = sz / 2;
                                                               S q = c; for (int &x: q) x = sub(0, x); q[0]=1;
   ntt(X.data(), sz, true); });
                                                               S p = Mul(a, q); p.resize(sz); q.resize(sz);
                                                               for (int r; r = (k & 1), k; k >>= 1) {
S Dx(S A) {
                                                                fill(d + all(p), 0); fill(d + 1 + all(q), 0);
fi(1, A.size()) A[i - 1] = mul(i, A[i]);
return A.empty() ? A : (A.pop_back(), A);
                                                                ntt(p.data(), sz); ntt(q.data(), sz);
                                                                for (size_t i = 0; i < sz; i++)</pre>
                                                                 p[i] = mul(p[i], q[(i + o) & (sz - 1)]);
S Sx(S A) {
                                                                for (size_t i = 0, j = 0; j < sz; i++, j++)</pre>
A.insert(A.begin(), 0);
                                                                q[i] = q[j] = mul(q[i], q[j]);
ntt(p.data(), sz, true); ntt(q.data(), sz, true);
fi(1, A.size()) A[i] = mul(modinv(int(i)), A[i]);
return A:
                                                                for (size_t i = 0; i < d; i++) p[i] = p[i << 1 | r];</pre>
                                                               for (size_t i = 0; i <= d; i++) q[i] = q[i << 1];
} // Bostan-Mori</pre>
S Ln(const S &A) { // coef[0] == 1; res[0] == 0
auto B = Sx(Mul(Dx(A), Inv(A), bit_ceil(A.size()*2)));
return B.resize(A.size()), B;
                                                               return mul(p[0], modinv(q[0]));
                                                              } // a_n = \sum c_j a_(n-j), c_0 is not used
5.16 Partition Number [9bb845]
S Exp(const S &v) { // coef[0] == 0; res[0] == 1
                                                              ans[0] = tmp[0] = 1;

for (int i = 1; i * i <= n; i++) {
return Newton(v, 1,
  [](S &X, S &A, int sz) {
                                                               for (int rep = 0; rep < 2; rep++)
for (int j = i; j <= n - i * i; j++)</pre>
   auto Y = X; Y.resize(sz / 2); Y = Ln(Y);
   fi(0, Y.size()) Y[i] = sub(A[i], Y[i]);
   Y[0] = add(Y[0], 1); X = Mul(X, Y, sz); );
                                                                 modadd(tmp[j], tmp[j-i]);
                                                               for (int j = i * i; j <= n; j++)
modadd(ans[j], tmp[j - i * i]);</pre>
S Pow(S a, lld M) { // period mod*(mod-1)
assert(!a.empty() && a[0] != 0);
                                                                     Pi Count 17158631
const auto imul = [&a](int s) {
                                                              struct S { int rough; lld large; int id; };
  for (int &x: a) x = mul(x, s); }; int c = a[0];
imul(modinv(c)); a = Ln(a); imul(int(M % mod));
                                                              lld PrimeCount(lld n) { // n \sim 10^{13} \Rightarrow < 1s
a = Exp(a); imul(modpow(c, int(M % (mod - 1))));
                                                               if (n <= 1) return 0;
 return a; // mod x^N where N=a.size()
                                                               const int v = static_cast<int>(sqrtl(n)); int pc = 0;
                                                               vector<int> smalls(v + 1), skip(v + 1); vector<S> z;
```

for (llu a2; t--; a = a2) {

```
for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
                                                                     a2 = mmul(a, a, x);
 for (int i : views::iota(0, (v + 1) / 2))
z.emplace_back(2*i+1, (n / (2*i+1) + 1) / 2, i);
                                                                    if (a2 == 1 && a != 1 && a != x - 1) return true;
                                                                   }
 for (int p = 3; p <= v; ++p)
if (smalls[p] > smalls[p - 1]) {
                                                                   return a != 1;
  const int q = p * p; ++pc;
                                                                  if (x <= 2 || ~x & 1) return x == 2;
  if (1LL * q * q > n) break;
                                                                  int t = countr_zero(x-1); llu odd = (x-1) >> t;
  skip[p] = 1;
                                                                  for (llu m:
  for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
                                                                   {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
  int ns = 0;
                                                                   if (m % x != 0 && witn(mpow(m % x, odd, x), t))
  for (auto e : z) if (!skip[e.rough]) {
                                                                    return false;
  lld d = 1LL * e.rough * p;
                                                                  return true;
   e.large += pc - (d <= v ? z[smalls[d] - pc].large :</pre>
                                                                 } // test @ luogu 143 & yosupo judge, ~1700ms for Q=1e5 \,
                                                                  // if use montgomery, ~250ms for Q=1e5
    smalls[n / d]);
   e.id = ns; z[ns++] = e;
                                                                 5.20 Pollard Rho [57ad88]
  }
                                                                 // does not work when n is prime or n == 1
  for (int j = v / p; j >= p; --j) {
                                                                 // return any non-trivial factor
                                                                 llu pollard_rho(llu n) {
   int c = smalls[j] - pc, e = min(j * p + p, v + 1);
                                                                  static mt19937_64 rnd(120821011);
   for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
                                                                  if (!(n & 1)) return 2;
  }
                                                                  llu y = 2, z = y, c = rnd() % n, p = 1, i = 0, t;
                                                                  auto f = [&](llu x) {
 lld ans = z[0].large; z.erase(z.begin());
                                                                   return madd(mmul(x, x, n), c, n); };
 for (auto &[rough, large, k] : z) {
                                                                  do {
  const lld m = n / rough; --k;
                                                                   p = mmul(msub(z = f(f(z)), y = f(y), n), p, n);
if (++i &= 63) if (i == (i & -i)) t = gcd(p, n);
  ans -= large - (pc + k);
 for (auto [p, ]
                  _, l] : z)
                                                                  } while (t == 1);
   if (l >= k || p * p > m) break;
                                                                  return t == n ? pollard_rho(n) : t;
   else ans += smalls[m / p] - (pc + l);
                                                                 } // test @ yosupo judge, ~270ms for Q=100
// if use montgomery, ~70ms for Q=100
 return ans;
                                                                 5.21 Montgomery [648fb3]
} // test @ yosupo library checker w/ n=1e11, 68ms
5.18 Min 25 Sieve [3695ef]
                                                                 struct Mont { // Montgomery multiplication
constexpr static int W = 64, L = 6;
template <typename U, typename V> struct min25 {
                                                                  llu mod, R1, R2, xinv;
 lld n; int sq;
                                                                  void set_mod(llu _mod) {
 vector<U> Ss, Sl, Spre; vector<V> Rs, Rl;
                                                                   mod = _mod; assert(mod & 1); xinv = 1;
 Sieve sv; vector<lld> quo;
                                                                   for (int j = 0; j < L; j++) xinv *= 2 - xinv * mod;</pre>
 U &S(lld d) { return d < sq ? Ss[d] : Sl[n / d]; }
                                                                   assert(xinv * mod == 1);
 V &R(lld d) { return d < sq ? Rs[d] : Rl[n / d]; }</pre>
                                                                   const u128 R = (u128(1) << W) % mod;</pre>
 min25(lld n_{-}) : n(n_{-}), sq((int)sqrt(n) + 1),
                                                                   R1 = llu(R); R2 = llu(R*R \% mod);
 Ss(sq), Sl(sq), Spre(sq), Rs(sq), Rl(sq), sv(sq) {
for (lld i = 1, Q; i <= n; i = n / Q + 1)</pre>
                                                                  llu redc(llu a, llu b) const {
   quo.push_back(Q = n / i);
                                                                   u128 T = u128(a) * b, m = -llu(T) * xinv;
                                                                   T += m * mod; T >>= W;
U F_prime(auto &&f, auto &&F) {
  for (lld p : sv.primes) Spre[p] = f(p);
                                                                   return llu(T >= mod ? T - mod : T);
  for (int i = 1; i < sq; i++) Spre[i] += Spre[i - 1];
for (lld i : quo) S(i) = F(i) - F(1);</pre>
                                                                  llu from(llu x) const {
                                                                  assert(x < mod); return redc(x, R2); }
llu get(llu a) const { return redc(a, 1); }</pre>
  for (lld p : sv.primes)
   for (lld i : quo) {
                                                                  llu one() const { return R1; }
    if (p * p > i) break;
                                                                 } mont;
    S(i) = f(p) * (S(i / p) - Spre[p - 1]);
                                                                  // a * b % mod == get(redc(from(a), from(b)))
                                                                 5.22 Berlekamp Massey [a94d00]
  return S(n);
                                                                 template <typename T>
 } // F_prime: \sum _ {p is prime, p <= n} f(p)
V F_comp(auto &&g) {</pre>
                                                                 vector<T> BerlekampMassey(const vector<T> &output) {
                                                                  vector<T> d(output.size() + 1), me, he;
  for (lld i : quo) R(i) = V(S(i));
                                                                  for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
  for (lld p : sv.primes | views::reverse)
                                                                   for (size_t j = 0; j < me.size(); ++j)
d[i] += output[i - j - 2] * me[j];</pre>
   for (lld i : quo) {
    if (p * p > i) break;
                                                                   if ((d[i] -= output[i - 1]) == 0) continue;
    lld prod = p;
                                                                   if (me.empty()) {
    for (int c = 1; prod * p <= i; ++c, prod *= p) {</pre>
                                                                    me.resize(f = i);
     R(i) += g(p, c) * (R(i / prod) - V(Spre[p]));
                                                                    continue;
     R(i) += g(p, c + 1);
                                                                   vector<T> o(i - f - 1);
                                                                   T k = -d[i] / d[f]; o.push_back(-k);
  return R(n);
                                                                   for (T x : he) o.push_back(x * k);
 }; // O(n^{3/4} / log n)
                                                                   if (o.size() < me.size()) o.resize(me.size());</pre>
                                                                   for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
/* U, V 都是環,記 h: U -> V 代表 U 轉型成 V 的函數。
                                                                   if (i-f+he.size() >= me.size()) he = me, f = i;
要求 h(x + y) = h(x) + h(y); f: lld -> U 是完全積性;
g 是積性函數且 h(f(p)) = g(p) 對於質數 p。
                                                                   me = o;
呼叫 F_comp 前需要先呼叫 F_prime 得到 S(i)。
S(i), R(i) 是 F_prime 和 F_comp 在 n/k 點的値。 F(i) = \sum_{i=1}^{n} f(i) 和 f(i) 需要快速求値。
                                                                  return me;
g(p, c) := g(pow(p, c)) 需要快速求值。
                                                                 5.23 Gauss Elimination [fa0977]
例如若 g(p) 是度數 d 的多項式則可以構造 f(p) 是維護
                                                                 using VI = vector<int>; // be careful if A.empty()
using VVI = vector<VI>; // ensure that 0 <= x < mod</pre>
pow(p, c) 的 (d+1)-tuple */
5.19 Miller Rabin [fbd812]
                                                                 pair<VI, VVI> gauss(VVI A, VI b) { // solve Ax=b
bool isprime(llu x) {
                                                                  const int N = (int)A.size(), M = (int)A[0].size();
 auto witn = [&](llu a, int t) {
                                                                  vector<int> depv, free(M, true); int rk = 0;
                                                                  for (int i = 0; i < M; i++) {</pre>
```

```
int p = -1;
                                                                        int s = -1;
                                                                        for (int i = 0; i <= n; ++i) {
  for (int j = rk; j < N; j++)</pre>
   if (p == -1 || abs(A[j][i]) > abs(A[p][i]))
                                                                         if (!z && q[i] == -1) continue;
  p = j;
if (p == -1 || A[p][i] == 0) continue;
                                                                         if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
  swap(A[p], A[rk]); swap(b[p], b[rk]);
                                                                        if (s == -1 || d[x][s] > -eps) return true;
  const int inv = modinv(A[rk][i]);
                                                                        int r = -1;
  for (int &x : A[rk]) x = mul(x, inv);
                                                                        for (int i = 0; i < m; ++i) {</pre>
  b[rk] = mul(b[rk], inv);
for (int j = 0; j < N; j++) if (j != rk) {
  int z = A[j][i];</pre>
                                                                         if (d[i][s] < eps) continue;</pre>
                                                                         if (r == -1 |
                                                                          d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
   for (int k = 0; k < M; k++)
    A[j][k] = sub(A[j][k], mul(z, A[rk][k]));
                                                                        if (r == -1) return false;
   b[j] = sub(b[j], mul(z, b[rk]));
                                                                       pivot(r, s);
  depv.push_back(i); free[i] = false; ++rk;
                                                                     VD solve(const VVD &a, const VD &b, const VD &c) {
                                                                      m = (int)b.size(), n = (int)c.size();
 for (int i = rk; i < N; i++)</pre>
                                                                      d = VVD(m + 2, VD(n + 2));
  if (b[i] != 0) return {{}}, {{}}}; // not consistent
 VI x(M); VVI h;
                                                                      for (int i = 0; i < m; ++i)</pre>
                                                                       for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
 for (int i = 0; i < rk; i++) x[depv[i]] = b[i];</pre>
 for (int i = 0; i < M; i++) if (free[i]) {</pre>
                                                                      p.resize(m), q.resize(n + 1);
                                                                       for (int i = 0; i < m; ++i)</pre>
  h.emplace_back(M); h.back()[i] = 1;
                                                                       p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
  for (int j = 0; j < rk; j++)</pre>
   h.back()[depv[j]] = sub(0, A[j][i]);
                                                                      for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];</pre>
                                                                      q[n] = -1, d[m + 1][n] = 1;
                                                                      int r = 0;
 return {x, h}; // solution = x + span(h[i])
                                                                      for (int i = 1; i < m; ++i)</pre>
5.24
        CharPoly [cd559d]
                                                                        if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
#define rep(x, y, z) for (int x=y; x<z; x++)
using VI = vector<int>; using VVI = vector<VI>;
                                                                      if (d[r][n + 1] < -eps) {
                                                                       pivot(r, n);
void Hessenberg(VVI &H, int N) {
                                                                        if (!phase(1) || d[m + 1][n + 1] < -eps)</pre>
 for (int i = 0; i < N - 2; ++i) {
                                                                         return VD(n, -inf);
  for (int j = i + 1; j < N; ++j) if (H[j][i]) {
                                                                        for (int i = 0; i < m; ++i) if (p[i] == -1) {
                                                                         int s = min_element(d[i].begin(), d[i].end() - 1)
   rep(k, i, N) swap(H[i+1][k], H[j][k]);
                                                                              - d[i].begin();
   rep(k, 0, N) swap(H[k][i+1], H[k][j]);
   break;
                                                                         pivot(i, s);
                                                                       }
  if (!H[i + 1][i]) continue;
  for (int j = i + 2; j < N; ++j) {
                                                                      if (!phase(0)) return VD(n, inf);
   int co = mul(modinv(H[i + 1][i]), H[j][i]);
                                                                      for (int i = 0; i < m; ++i)</pre>
   rep(k, i, N) subeq(H[j][k], mul(H[i+1][k], co));
   rep(k, 0, N) addeq(H[k][i+1], mul(H[k][j], co));
                                                                       if (p[i] < n) x[p[i]] = d[i][n + 1];</pre>
  }
                                                                      return x;
                                                                     }} // use double instead of long double if possible
}
                                                                     5.26 Simplex Construction
VI CharacteristicPoly(VVI A) {
                                                                     Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j for all
 int N = (int)A.size(); Hessenberg(A, N);
                                                                     1 \le j \le m and x_i \ge 0 for all 1 \le i \le n.
 VVI P(N + 1, VI(N + 1)); P[0][0] = 1;
                                                                     1. In case of minimization, let c_i^\prime = -c_i
 for (int i = 1; i <= N; ++i) {
                                                                     \begin{array}{ll} \textbf{2.} & \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j \\ \textbf{3.} & \sum_{1 < i < n} A_{ji} x_i = b_j \rightarrow \mathsf{add} \leq \mathsf{and} \geq. \end{array}
  rep(j, 0, i+1) P[i][j] = j ? P[i-1][j-1] : 0;
  for (int j = i - 1, val = 1; j >= 0; --j) {
  int co = mul(val, A[j][i - 1]);
                                                                     4. If x_i has no lower bound, replace x_i with x_i - x_i'
                                                                     5.27 Adaptive Simpson [b8cef9]
   rep(k, 0, j+1) subeq(P[i][k], mul(P[j][k], co));
                                                                     llf integrate(auto &&f, llf L, llf R) {
   if (j) val = mul(val, A[j][j - 1]);
                                                                      auto simp = [&](llf l, llf r) {
  }
                                                                       llf m = (l + r) / 2;
                                                                        return (f(l) + f(r) + 4.0 * f(m)) * (r - l) / 6.0;
 if (N & 1) for (int &x: P[N]) x = sub(0, x);
 return P[N]; // test: 2021 PTZ Korea K
                                                                      auto F = [&](auto Y, llf l, llf r, llf v, llf eps) {
    llf m = (l+r)/2, vl = simp(l, m), vr = simp(m, r);
5.25 Simplex [c9c93b]
                                                                        if (abs(vl + vr - v) \le 15 * eps)
namespace simplex {
                                                                        return vl + vr + (vl + vr - v) / 15.0;
// maximize c^Tx under Ax \le B and x \ge 0
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
                                                                        return Y(Y, l, m, vl, eps / 2.0) +
                                                                                Y(Y, m, r, vr, eps / 2.0);
using VD = vector<llf>;
                                                                      return F(F, L, R, simp(L, R), 1e-6);
using VVD = vector<vector<llf>>;
const llf eps = 1e-9, inf = 1e+9;
                                                                     5.28 Poly Roots* [235182]
int n, m; VVD d; vector<int> p, q;
                                                                     VD polyRoots(VD p, llf xmin, llf xmax) {
void pivot(int r, int s) {
 llf inv = 1.0 / d[r][s];
                                                                      if (p.size() == 2) return {-p[0]/p[1]};
 for (int i = 0; i < m + 2; ++i)</pre>
                                                                      VD d = polyRoots(derivative(p), xmin, xmax), ret;
  for (int j = 0; j < n + 2; ++j)</pre>
                                                                      d.pb(xmin-1); d.pb(xmax+1); sort(all(d));
                                                                      for (size_t i = 0; i + 1 < d.size(); i++) {
    llf l = d[i], h = d[i+1]; bool s = eval(p, l) > 0;
   if (i != r && j != s)
    d[i][j] = d[r][j] * d[i][s] * inv;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
                                                                        if (s ^ (eval(p, h) > 0)) {
                                                                        for (int _ = 0; _ < 60; _++) {
    llf m = (l + h) / 2, f = eval(p, m);
 d[r][s] = inv; swap(p[r], q[s]);
                                                                          ((f \le 0) \land s ? l : h) = m;
bool phase(int z) {
 int x = m + z;
                                                                         ret.push_back((l + h) / 2);
 while (true) {
```

```
while (cross(e, p[(pos + 1) % n] - p[i]) >
                                                                  cross(e, p[pos] - p[i]))
 return ret:
                                                                pos = (pos + 1) % n;
                                                               for (int j: {i, (i + 1) % n})
ans = max(ans, norm(p[pos] - p[j]));
5.29 Golden Ratio Search [376bcb]
llf gss(llf a, llf b, auto &&f) {
                                                             } // tested @ AOJ CGL_4_B
 llf r = (sqrt(5)-1)/2, eps = 1e-7;
                                                              6.4 MinMax Enclosing Rect [e4470c]
 llf x1 = b - r*(b-a), x2 = a + r*(b-a);
 llf f1 = f(x1), f2 = f(x2);
                                                             // from 8BQube, plz ensure p is strict convex hull
 while (b-a > eps)
                                                             const llf INF = 1e18, qi = acos(-1) / 2 * 3;
                                                             pair<llf, llf> solve(const vector<P> &p) {
    llf mx = 0, mn = INF; int n = (int)p.size();
    for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {</pre>
  if (f1 < f2) { //change to > to find maximum
   b = x2; x2 = x1; f2 = f1;
   x1 = b - r*(b-a); f1 = f(x1);
  } else {
                                                             #define Z(v) (p[(v) % n] - p[i])
   a = x1; x1 = x2; f1 = f2;
                                                               P = Z(i + 1);
   x2 = a + r*(b-a); f2 = f(x2);
                                                                while (cross(e, Z(u + 1)) > cross(e, Z(u))) ++u;
 }
                                                                while (dot(e, Z(r + 1)) > dot(e, Z(r))) ++r;
 return a;
                                                                if (!i) l = r + 1;
                                                                while (dot(e, Z(l + 1)) < dot(e, Z(l))) ++l;</pre>
                                                                P D = p[r % n] - p[l % n];
     Geometry
                                                               llf H = cross(e, Z(u)) / llf(norm(e));
6.1 Basic Geometry [1d2d70]
                                                                mn = min(mn, dot(e, D) * H);
#define IM imag
                                                                llf B = sqrt(norm(D)) * sqrt(norm(Z(u)));
#define RE real
                                                               llf deg = (qi - acos(dot(D, Z(u)) / B)) / 2;
using lld = int64_t;
using llf = long double;
                                                               mx = max(mx, B * sin(deg) * sin(deg));
using PT = complex<lld>;
                                                              return {mn, mx};
using PF = complex<llf>;
                                                             } // test @ UVA 819
using P = PT;
                                                              6.5 Minkowski Sum [602806]
llf abs(P p) { return sqrtl(norm(p)); }
PF toPF(PT p) { return PF{RE(p), IM(p)}; }
int sgn(lld x) { return (x > 0) - (x < 0); }</pre>
                                                             // A, B are strict convex hull rotate to min by (X, Y)
                                                             vector<P> Minkowski(vector<P> A, vector<P> B) {
lld dot(P a, P b) { return RE(conj(a) * b); }
                                                               const int N = (int)A.size(), M = (int)B.size();
                                                              vector<P> sa(N), sb(M), C(N + M + 1);
for (int i = 0; i < N; i++) sa[i] = A[(i+1)%N]-A[i];</pre>
lld cross(P a, P b) { return IM(conj(a) * b); }
int ori(P a, P b, P c) {
                                                               for (int i = 0; i < M; i++) sb[i] = B[(i+1)%M]-B[i];</pre>
return sgn(cross(b - a, c - a));
                                                               C[0] = A[0] + B[0];
                                                               for (int i = 0, j = 0; i < N || j < M; ) {</pre>
int quad(P p) {
                                                               P = (j>=M \mid | (i<N && cross(sa[i], sb[j])>=0))
return (IM(p) == 0) // use sgn for PF
                                                                ? sa[i++] : sb[j++];
  ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
                                                               C[i + j] = e;
int argCmp(P a, P b) {
                                                               partial_sum(all(C), C.begin()); C.pop_back();
 // returns 0/+-1, starts from theta = -PI
 int qa = quad(a), qb = quad(b);
                                                               return convex_hull(C); // just to remove colinear
                                                             \} // be careful if min(|A|,|B|) \le 2
 if (qa != qb) return sgn(qa - qb);
 return sgn(cross(b, a));
                                                              6.6 Segment Intersection [60d016]
                                                              struct Seg { // closed segment
                                                              P st, dir; // represent st + t*dir for 0 <= t <= 1 Seg(P s, P e) : st(s), dir(e - s) {}
P rot90(P p) { return P{-IM(p), RE(p)}; }
template <typename V> llf area(const V & pt) {
 lld ret = 0;
                                                               static bool valid(lld p, lld q) {
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
                                                                // is there t s.t. 0 <= t <= 1 && qt == p ?
 ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
                                                                if (q < 0) q = -q, p = -p;
 return ret / 2.0;
                                                               return 0 <= p && p <= q;
template <typename V> PF center(const V & pt) {
                                                               vector<P> ends() const { return { st, st + dir }; }
 P ret = 0; lld A = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++) {</pre>
                                                             template <typename T> bool isInter(T A, P p) {
  lld cur = cross(pt[i] - pt[0], pt[i+1] - pt[0]);
                                                               if (A.dir == P(0)) return p == A.st; // BE CAREFUL
  ret += (pt[i] + pt[i + 1] + pt[0]) * cur; A += cur;
                                                               return cross(p - A.st, A.dir) == 0 &&
                                                                T::valid(dot(p - A.st, A.dir), norm(A.dir));
return toPF(ret) / llf(A * 3);
                                                             template <typename U, typename V>
PF project(PF p, PF q) { // p onto q
                                                             bool isInter(U A, V B) {
return dot(p, q) * q / dot(q, q); // dot<llf>
                                                               if (cross(A.dir, B.dir) == 0) { // BE CAREFUL
                                                               bool res = false;
6.2 2D Convex Hull [ecba37]
                                                                for (P p: A.ends()) res |= isInter(B, p);
                                                                for (P p: B.ends()) res |= isInter(A, p);
// from NaCl, counterclockwise, be careful of n<=2
vector<P> convex_hull(vector<P> v) { // n==0 will RE
                                                                return res;
 sort(all(v)); // by X then Y
                                                               P D = B.st - A.st; lld C = cross(A.dir, B.dir);
 if (v[0] == v.back()) return {v[0]};
                                                               return U::valid(cross(D, B.dir), C) &&
 int t = 0, s = 1; vector<P> h(v.size() + 1);
                                                                V::valid(cross(D, A.dir), C);
 for (int _ = 2; _--; s = t--, reverse(all(v)))
  for (P p : v) {
   while (t>s && ori(p, h[t-1], h[t-2]) >= 0) t--;
                                                              6.7 Halfplane Intersection [f2bd8f]
   h[t++] = p;
                                                             struct Line {
                                                              P st, ed, dir;
 return h.resize(t), h;
                                                               Line (P s, P e) : st(s), ed(e), dir(e - s) {}
                                                             }; using LN = const Line &;
                                                             PF intersect(LN A, LN B) {
6.3 2D Farthest Pair [8b5844]
                                                              llf t = cross(B.st - A.st, B.dir) /
// p is CCW convex hull w/o colinear points
int n = (int)p.size(), pos = 1; lld ans = 0;
                                                               llf(cross(A.dir, B.dir));
for (int i = 0; i < n; i++) {</pre>
                                                               return toPF(A.st) + toPF(A.dir) * t; // C^3 / C^2
P = p[(i + 1) \% n] - p[i];
```

```
bool cov(LN l, LN A, LN B) {
                                                                sort(all(e));
                                                                vector<int> ord(n), pos(n);
 i128 u = cross(B.st-A.st, B.dir);
 i128 v = cross(A.dir, B.dir);
                                                                iota(all(ord), 0);
                                                                sort(all(ord), [&p](int i, int j) {
  return cmpxy(p[i], p[j]); });
 // ori(l.st, l.ed, A.st + A.dir*(u/v)) <= 0?
 i128 x = RE(A.dir) * u + RE(A.st - l.st) * v;
 i128 y = IM(A.dir) * u + IM(A.st - l.st) * v;
                                                                for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
 return sgn(x*IM(l.dir) - y*RE(l.dir)) * sgn(v) >= 0;
                                                                const auto makeReverse = [](auto &v) {
} // x, y are C^3, also sgn<i128> is needed
                                                                 sort(all(v)); v.erase(unique(all(v)), v.end());
                                                                 vector<pair<int,int>> segs;
bool operator<(LN a, LN b) {</pre>
                                                                 for (size_t i = 0, j = 0; i < v.size(); i = j) {
   for (; j < v.size() && v[j] - v[i] <= j - i; j++);</pre>
 if (int c = argCmp(a.dir, b.dir)) return c == -1;
 return ori(a.st, a.ed, b.st) < 0;</pre>
                                                                  segs.emplace_back(v[i], v[j-1] + 1 + 1);
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
  the half plane is the LHS when going from st to ed
                                                                 return segs;
llf HPI(vector<Line> &q) {
                                                                };
 sort(q.begin(), q.end());
                                                                for (size_t i = 0, j = 0; i < e.size(); i = j) {</pre>
 int n = (int)q.size(), l = 0, r = -1;
                                                                 /* do here */
 for (int i = 0; i < n; i++) {
                                                                 vector<size_t> tmp;
  if (i && !argCmp(q[i].dir, q[i-1].dir)) continue;
                                                                 for (; j < e.size() && !(e[i] < e[j]); j++)</pre>
  while (l < r && cov(q[i], q[r-1], q[r])) --r;
                                                                  tmp.push_back(min(pos[e[j].u], pos[e[j].v]));
  while (l < r && cov(q[i], q[l], q[l+1])) ++l;</pre>
                                                                 for (auto [l, r] : makeReverse(tmp)) {
  q[++r] = q[i];
                                                                  reverse(ord.begin() + l, ord.begin() + r);
                                                                  for (int t = l; t < r; t++) pos[ord[t]] = t;</pre>
 while (l < r && cov(q[l], q[r-1], q[r])) --r;</pre>
 while (l < r && cov(q[r], q[l], q[l+1])) ++l;</pre>
                                                               }
n = r - l + 1; // q[l .. r] are the lines
if (n \le 2 \mid | !argCmp(q[l].dir, q[r].dir)) return 0;
                                                               6.11
                                                                     Hull Cut* [277def]
 vector<PF> pt(n);
                                                               vector<P> cut(const vector<P> &p, P s, P e) {
 for (int i = 0; i < n; i++)</pre>
                                                                vector<P> res;
  pt[i] = intersect(q[i+l], q[(i+1)%n+l]);
                                                                for (size_t i = 0; i < p.size(); i++) {</pre>
 return area(pt);
                                                                 P cur = p[i], prv = i ? p[i-1] : p.back();
} // test @ 2020 Nordic NCPC : BigBrother
                                                                 bool side = ori(s, e, cur) < 0;</pre>
6.8 HPI alternative form* [c77a2f]
                                                                 if (side != (ori(s, e, prv) < 0))
                                                                  res.push_back(intersect({s, e}, {cur, prv}));
struct Line {
                                                                 if (side) res.push_back(cur);
 lld a, b, c; // ax + by + c \le 0
 P dir() const { return P(a, b); }
                                                                } // P is complex<llf>
                                                              return res; // hull intersection with halfplane
} // left of the line s -> e
}; using LN = const Line &;
PF intersect(LN A, LN B) {
 P x(A.a, B.a), y(A.b, B.b), z(A.c, B.c);
                                                               6.12 Point In Hull [13edeb]
 llf D = cross(x, y);
                                                              bool isAnti(P a, P b) {
 return PF(cross(y, z) / D, cross(z, x) / D);
                                                               return cross(a, b) == 0 && dot(a, b) <= 0; }</pre>
                                                               bool PIH(const vector<P> &h, P z, bool strict = true) {
bool cov(LN l, LN A, LN B) {
                                                                int n = (int)h.size(), a = 1, b = n - 1, r = !strict;
 P x(A.a, B.a), y(A.b, B.b), z(A.c, B.c);
                                                                if (n < 3) return r && isAnti(h[0] - z, h[n-1] - z);</pre>
 i128 c = cross(x,y), a = cross(y,z), b = cross(z,x);
                                                                if (ori(h[0],h[a],h[b]) > 0) swap(a, b);
 return sgn(a * l.a + b * l.b + c * l.c) * sgn(c) > 0;
                                                                if (ori(h[0],h[a],z) >= r || ori(h[0],h[b],z) <= -r)</pre>
                                                                 return false;
bool operator<(LN a, LN b) {</pre>
                                                                while (abs(a - b) > 1) {
 if (int c = argCmp(a.dir(), b.dir())) return c == -1;
                                                                 int c = (a + b) / 2;
 return a.c > b.c;
                                                                 (ori(h[0], h[c], z) > 0 ? b : a) = c;
      SegmentDist (Sausage) [9d8603]
                                                                return ori(h[a], h[b], z) < r;
// be careful of abs<complex<int>> (replace _abs below)
                                                                     Point In Polygon [037c52]
llf PointSegDist(P A, Seg B) {
 if (B.dir == P(0)) return _abs(A - B.st);
                                                              bool PIP(const vector\langle P \rangle &p, P z, bool strict = true) {
 if (sgn(dot(A - B.st, B.dir)) *
                                                                int cnt = 0, n = (int)p.size();
                                                                for (int i = 0; i < n; i++) {</pre>
   sgn(dot(A - B.ed, B.dir)) <= 0)</pre>
  return abs(cross(A - B.st, B.dir)) / _abs(B.dir);
                                                                P A = p[i], B = p[(i + 1) \% n];
                                                                if (isInter(Seg(A, B), z)) return !strict;
auto zy = IM(z), Ay = IM(A), By = IM(B);
 return min(_abs(A - B.st), _abs(A - B.ed));
llf SegSegDist(const Seg &s1, const Seg &s2) {
                                                                cnt ^= ((zy<Ay) - (zy<By)) * ori(z, A, B) > 0;
 if (isInter(s1, s2)) return 0;
                                                                }
 return min({
                                                                return cnt;
   PointSegDist(s1.st, s2),
   PointSegDist(s1.ed, s2),
                                                               6.14 Point In Polygon (Fast) [2cd3d6]
   PointSegDist(s2.st, s1),
   PointSegDist(s2.ed, s1) });
                                                               vector<int> PIPfast(vector<P> p, vector<P> q) {
} // test @ QOJ2444 / PTZ19 Summer.D3
                                                               const int N = int(p.size()), Q = int(q.size());
6.10 Rotating Sweep Line [8aff27]
                                                                vector<pair<P, int>> evt; vector<Seg> edge;
struct Event {
                                                                for (int i = 0; i < N; i++) {</pre>
                                                                 int a = i, b = (i + 1) % N;
 Pd; int u, v;
 bool operator<(const Event &b) const {</pre>
                                                                P A = p[a], B = p[b];
                                                                 assert (A < B || B < A); // std::operator<</pre>
  return sgn(cross(d, b.d)) > 0; }
                                                                 if (B < A) swap(A, B);
P makePositive(P z) { return cmpxy(z, 0) ? -z : z; }
                                                                 evt.emplace_back(A, i); evt.emplace_back(B, ~i);
void rotatingSweepLine(const vector<P> &p) {
                                                                edge.emplace_back(A, B);
 const int n = int(p.size());
 vector<Event> e; e.reserve(n * (n - 1) / 2);
                                                                for (int i = 0; i < Q; i++)</pre>
 for (int i = 0; i < n; i++)
for (int j = i + 1; j < n; j++)</pre>
                                                                evt.emplace_back(q[i], i + N);
                                                                sort(all(evt));
                                                                auto vtx = p; sort(all(vtx));
   e.emplace_back(makePositive(p[i] - p[j]), i, j);
```

```
auto eval = [](const Seg &a, lld x) -> llf {
                                                              llf phi = acos((a.r * a.r + d2 - b.r * b.r) /
                                                                 (2 * a.r * dis)); // is acos_safe needed?
  if (RE(a.dir) == 0) {
   assert (x == RE(a.st));
                                                              llf L = FMOD(theta - phi), R = FMOD(theta + phi);
   return IM(a.st) + llf(IM(a.dir)) / 2;
                                                              return { L, R };
  llf t = (x - RE(a.st)) / llf(RE(a.dir));
                                                             vector<PF> intersectPoint(Cir a, Cir b) {
  return IM(a.st) + IM(a.dir) * t;
                                                              llf d = abs(a.o - b.o);
                                                              if (d > b.r+a.r || d < abs(b.r-a.r)) return {};</pre>
 lld cur_x = 0;
                                                              llf dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
 auto cmp = [&](const Seg &a, const Seg &b) -> bool {
  if (int s = sgn(eval(a, cur_x) - eval(b, cur_x)))
                                                              PF dir = (a.o - b.o) / d;
                                                              PF u = dir * d1 + b.o;
   return s == -1; // be careful: sgn<llf>, sgn<lld>
                                                              PF v = rot90(dir) * sqrt(max(0.0L, b.r*b.r-d1*d1));
                                                              return {u + v, u - v};
  int s = sgn(cross(b.dir, a.dir));
  if (cur_x != RE(a.st) && cur_x != RE(b.st)) s *= -1;
                                                             } // test @ AOJ CGL probs
  return s == -1;
                                                             6.18 Circle Common Tangent [d97f1c]
 };
                                                             // be careful of tangent / exact same circle
 namespace pbds = __gnu_pbds;
                                                             // sign1 = 1 for outer tang, -1 for inter tang
 pbds::tree<Seg, int, decltype(cmp),</pre>
                                                             vector<Line> common_tan(const Cir &a, const Cir &b, int
  pbds::rb_tree_tag,
                                                                   sign1) {
  pbds::tree_order_statistics_node_update> st(cmp);
                                                              if (norm(a.o - b.o) < eps) return {};</pre>
 auto answer = [&](P ep) {
                                                              llf d = abs(a.o - b.o), c = (a.r - sign1 * b.r) / d;
  if (binary_search(all(vtx), ep))
                                                              PF v = (b.o - a.o) / d;
   return 1; // on vertex
                                                              if (c * c > 1) return {};
  Seg H(ep, ep); // ??
                                                              if (abs(c * c - 1) < eps) {
  auto it = st.lower_bound(H);
                                                               PF p = a.o + c * v * a.r;
  if (it != st.end() && isInter(it->first, ep))
                                                               return {Line(p, p + rot90(b.o - a.o))};
   return 1; // on edge
  if (it != st.begin() && isInter(prev(it)->first, ep))
                                                              vector<Line> ret; llf h = sqrt(max(0.0L, 1-c*c));
  return 1; // on edge
                                                              for (int sign2 : {1, -1}) {
  auto rk = st.order_of_key(H);
                                                               PF n = c * v + sign2 * h * rot90(v);
  return rk % 2 == 0 ? 0 : 2; // 0: outside, 2: inside
                                                               PF p1 = a.o + n * a.r;
 };
                                                               PF p2 = b.o + n * (b.r * sign1);
 vector<int> ans(Q);
                                                               ret.emplace_back(p1, p2);
 for (auto [ep, i] : evt) {
  cur_x = RE(ep);
                                                              return ret;
  if (i < 0) { // remove
  st.erase(edge[~i]);
                                                             6.19 Line-Circle Intersection [10786a]
 } else if (i < N) { // insert</pre>
                                                             vector<PF> LineCircleInter(PF p1, PF p2, PF o, llf r) {
   auto [it, succ] = st.insert({edge[i], i});
                                                              PF ft = p1 + project(o-p1, p2-p1), vec = p2-p1;
   assert(succ);
                                                              llf dis = abs(o - ft);
  } else ans[i - N] = answer(ep);
                                                              if (abs(dis - r) < eps) return {ft};</pre>
                                                              if (dis > r) return {};
 return ans;
                                                              vec = vec * sqrt(r * r - dis * dis) / abs(vec);
} // test @ AOJ CGL_3_C
                                                              return {ft + vec, ft - vec}; // sqrt_safe?
6.15 Cyclic Ternary Search [162adf]
int cyclic_ternary_search(int N, auto &&lt_) {
                                                             6.20 Poly-Circle Intersection [8e5]331
 auto lt = [&](int x, int y) {
                                                             // Divides into multiple triangle, and sum up
  return lt_(x % N, y % N); };
                                                             // from 8BQube, test by HDU2892 & AOJ CGL_7_H
int l = 0, r = N; bool up = lt(0, 1);
while (r - l > 1) {
                                                             llf _area(PF pa, PF pb, llf r) {
  if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
  int m = (l + r) / 2;
                                                              if (abs(pb) < eps) return 0;</pre>
  if (lt(m, 0) ? up : !lt(m, m+1)) r = m;
                                                              llf S, h, theta;
 else l = m;
                                                              llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
                                                              llf cB = dot(pb, pb-pa) / a / c, B = acos_safe(cB);
 return (lt(l, r) ? r : l) % N;
                                                              llf cC = dot(pa, pb) / a / b, C = acos_safe(cC);
} // find maximum; be careful if N == 0
                                                              if (a > r) {
6.16 Tangent of Points to Hull [8e1343]
                                                               S = (C / 2) * r * r; h = a * b * sin(C) / c;
pair<int, int> get_tangent(const vector<P> &v, P p) {
                                                               if (h < r && B < PI / 2)
 auto gao = [&](int s) {
                                                                 S = (acos\_safe(h/r)*r*r - h*sqrt\_safe(r*r-h*h));
  return cyclic_ternary_search(v.size(),
                                                              } else if (b > r) {
    [&](int x, int y) {
                                                               theta = PI - B - asin_safe(sin(B) / r * a);
     return ori(p, v[x], v[y]) == s; });
                                                               S = 0.5 * a*r*sin(theta) + (C-theta)/2 * r * r;
 }; // test @ codeforces.com/gym/101201/problem/E
                                                              } else
 return {gao(1), gao(-1)}; // (a,b):ori(p,v[a],v[b])<0
                                                               S = 0.5 * sin(C) * a * b;
} // plz ensure that point strictly out of hull
// if colinear, returns arbitrary point on line
                                                              return S;
6.17 Circle Class & Intersection [d5df51]
                                                             llf area_poly_circle(const vector<PF> &v, PF 0, llf r)
llf FMOD(llf x) {
                                                              llf S = 0;
 if (x < -PI) x += PI * 2;
                                                              for (size_t i = 0, N = v.size(); i < N; ++i)</pre>
 if (x > PI) x -= PI * 2;
                                                               S += _area(v[i] - 0, v[(i + 1) % N] - 0, r) *
return x;
                                                                  ori(0, v[i], v[(i + 1) % N]);
                                                              return abs(S);
struct Cir { PF o; llf r; };
// be carefule when tangent
                                                             6.21 Min Covering Circle [054ee0]
vector<llf> intersectAngle(Cir a, Cir b) {
PF dir = b.o - a.o; llf d2 = norm(dir);

if (norm(a.r - b.r) >= d2) { // norm(x) := |x|^2
                                                             Cir getCircum(P a, P b, P c){ // P = complex<llf>
P z1 = a - b, z2 = a - c; llf D = cross(z1, z2) * 2;
 if (a.r < b.r) return {-PI, PI}; // a in b</pre>
                                                              auto c1 = dot(a + b, z1), c2 = dot(a + c, z2);
                                                              P o = rot90(c2 * z1 - c1 * z2) / D;
  else return {}; // b in a
 } else if (norm(a.r + b.r) <= d2) return {};</pre>
                                                              return { o, abs(o - a) };
llf dis = abs(dir), theta = arg(dir);
```

```
Cir minCircleCover(vector<P> p) { // what if p.empty?
 Cir c = { 0, 0 }; shuffle(all(p), mt19937(114514));
 for (size_t i = 0; i < p.size(); i++) {</pre>
  if (abs(p[i] - c.o) <= c.r) continue;</pre>
  c = { p[i], 0 };
  for (size_t j = 0; j < i; j++) {</pre>
   if (abs(p[j] - c.o) <= c.r) continue;</pre>
   c.o = (p[i] + p[j]) / llf(2);
   c.r = abs(p[i] - c.o);
   for (size_t k = 0; k < j; k++) {</pre>
    if (abs(p[k] - c.o) <= c.r) continue;</pre>
    c = getCircum(p[i], p[j], p[k]);
   }
  }
 return c;
} // test @ TIOJ 1093 & luogu P1742
6.22 Circle Union [073c1c]
#define eb emplace_back
struct Teve { // test@SPOJ N=1000, 0.3~0.5s
PF p; llf a; int add; // point, ang, add
 Teve(PF x, llf y, int z) : p(x), a(y), add(z) {}
                                                                 struct P3 {
                                                                  lld x, y, z;
 bool operator<(Teve &b) const { return a < b.a; }</pre>
// strict: x = 0, otherwise x = -1
bool disjunct(Cir &a, Cir &b, int x)
{ return sgn(abs(a.o - b.o) - a.r - b.r) > x; }
bool contain(Cir &a, Cir &b, int x)
{ return sgn(a.r - b.r - abs(a.o - b.o)) > x; }
vector<llf> CircleUnion(vector<Cir> &c) {
 // area[i] : area covered by at least i circles
 int N = (int)c.size(); vector<llf> area(N + 1);
 vector<vector<int>> overlap(N, vector<int>(N));
 auto g = overlap; // use simple 2darray to speedup
 for (int i = 0; i < N; ++i)
  for (int j = 0; j < N; ++j) {
   /* c[j] is non-strictly in c[i]. */
   overlap[i][j] = i != j &&
    (sgn(c[i].r - c[j].r) > 0 ||
(sgn(c[i].r - c[j].r) == 0 && i < j)) &&
    contain(c[i], c[j], -1);
                                                                 struct Face {
 for (int i = 0; i < N; ++i)</pre>
                                                                  int a, b, c;
  for (int j = 0; j < N; ++j)</pre>
   g[i][j] = i != j && !(overlap[i][j] ||
     overlap[j][i] || disjunct(c[i], c[j], -1));
 for (int i = 0; i < N; ++i) {
  vector<Teve> eve; int cnt = 1;
  for (int j = 0; j < N; ++j) cnt += overlap[j][i];</pre>
  // if (cnt > 1) continue; (if only need area[1])
  for (int j = 0; j < N; ++j) if (g[i][j]) {</pre>
   auto IP = intersectPoint(c[i], c[j]);
   PF aa = IP[1], bb = IP[0];
   llf A = arg(aa - c[i].o), B = arg(bb - c[i].o);
   eve.eb(bb, B, 1); eve.eb(aa, A, -1);
if (B > A) ++cnt;
  if (eve.empty()) area[cnt] += PI*c[i].r*c[i].r;
   sort(eve.begin(), eve.end());
   eve.eb(eve[0]); eve.back().a += PI * 2;
for (size_t j = 0; j + 1 < eve.size(); j++) {</pre>
    cnt += eve[j].add;
    area[cnt] += cross(eve[j].p, eve[j+1].p) *.5;
    llf t = eve[j + 1].a - eve[j].a;
    area[cnt] += (t-sin(t)) * c[i].r * c[i].r *.5;
   }
  }
}
 return area;
6.23 Polygon Union [42e75b]
llf polyUnion(const vector<vector<P>> &p) {
vector<tuple<P, P, int>> seg;
for (int i = 0; i < ssize(p); i++)</pre>
  for (int j = 0, m = int(p[i].size()); j < m; j++)</pre>
   seg.emplace_back(p[i][j], p[i][(j + 1) % m], i);
 llf ret = 0; // area of p[i] must be non-negative
 for (auto [A, B, i] : seg) {
  vector<pair<llf, int>> evt{{0, 0}, {1, 0}};
for (auto [C, D, j] : seg) {
                                                                   now = next;
   int sc = ori(A, B, C), sd = ori(A, B, D);
```

```
if (sc != sd && i != j && min(sc, sd) < 0) {
    llf sa = cross(D-C, A-C), sb = cross(D-C, B-C);
    evt.emplace_back(sa / (sa - sb), sgn(sc - sd));
   } else if (!sc && !sd && j < i
     && sgn(dot(B - A, D - C)) > 0) {
    evt.emplace_back(real((C - A) / (B - A)), 1);
    evt.emplace_back(real((D - A) / (B - A)), -1);
  for (auto &[q, _] : evt) q = clamp<llf>(q, 0, 1);
sort(evt.begin(), evt.end());
  llf sum = 0, last = 0; int cnt = 0;
  for (auto [q, c] : evt) {
   if (!cnt) sum += q - last;
   cnt += c; last = q;
  ret += cross(A, B) * sum;
 return ret / 2;
6.24 3D Point [46b73b]
 P3 operator^(const P3 &b) const {
  return {y*b.z-b.y*z, z*b.x-b.z*x, x*b.y-b.x*y};
 //Azimuthal angle (longitude) to x-axis. \in [-pi, pi]
 llf phi() const { return atan2(y, x); }
 //Zenith angle (latitude) to the z-axis. \in [0, pi]
 llf theta() const { return atan2(sqrt(x*x+y*y),z); }
P3 ver(P3 a, P3 b, P3 c) { return (b - a) ^ (c - a); }
lld volume(P3 a, P3 b, P3 c, P3 d) {
return dot(ver(a, b, c), d - a);
P3 rotate_around(P3 p, llf angle, P3 axis) {
llf s = sin(angle), c = cos(angle);
 P3 u = normalize(axis);
 return u*dot(u, p)*(1-c) + p * c + cross(u, p)*s;
6.25 3D Convex Hull [01652a]
 Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
auto preprocess(const vector<P3> &pt) {
auto G = pt.begin();
 auto a = find_if(all(pt), [&](P3 z) {
  return z != *G; }) - G;
 auto b = find_if(all(pt), [&](P3 z) {
  return ver(*G, pt[a], z) != P3(0, 0, 0); }) - G;
 auto c = find_if(all(pt), [&](P3 z) {
  return volume(*G, pt[a], pt[b], z) != 0; }) - G;
 vector<size_t> id;
 for (size_t i = 0; i < pt.size(); i++)</pre>
 if (i != a && i != b && i != c) id.push_back(i);
 return tuple{a, b, c, id};
// return the faces with pt indexes
// all points coplanar case will WA
vector<Face> convex_hull_3D(const vector<P3> &pt) {
 const int n = int(pt.size());
 if (n <= 3) return {}; // be careful about edge case</pre>
 vector<Face> now;
 vector<vector<int>> z(n, vector<int>(n));
 auto [a, b, c, ord] = preprocess(pt);
 now.emplace_back(a, b, c); now.emplace_back(c, b, a);
 for (auto i : ord) {
  vector<Face> next;
  for (const auto &f : now) {
   lld v = volume(pt[f.a], pt[f.b], pt[f.c], pt[i]);
   if (v <= 0) next.push_back(f);</pre>
   z[f.a][f.b] = z[f.b][f.c] = z[f.c][f.a] = sgn(v);
  const auto F = [\&](int x, int y) \{
   if (z[x][y] > 0 && z[y][x] <= 0)
   next.emplace_back(x, y, i);
  for (const auto &f : now)
   F(f.a, f.b), F(f.b, f.c), F(f.c, f.a);
```

```
return now;
                                                                      link(E(X, 0), E(Y, 0));
                                                                      link(E(X, 1), A \rightarrow e[L(a)]); link(E(X, 2), B \rightarrow e[R(b)]);
// n^2 delaunay: facets with negative z normal of
                                                                      link(E(Y, 1), B\rightarrow e[L(b)]); link(E(Y, 2), A\rightarrow e[R(a)]);
// convexhull of (x, y, x^2 + y^2), use a pseudo-point // (0, 0, inf) to avoid degenerate case
                                                                      A\rightarrow ch = B\rightarrow ch = \{X, Y, nullptr\};
                                                                      flip(X, 1); flip(X, 2); flip(Y, 1); flip(Y, 2);
// test @ SPOJ CH3D
// llf area = 0, vol = 0; // surface area / volume
                                                                     void add_point(int p) {
                                                                      Tri *r = root;
// for (auto [a, b, c]: faces)
// area += abs(ver(p[a], p[b], p[c]))/2.0,
// vol += volume(P3(0, 0, 0), p[a], p[b], p[c])/6.0;
                                                                      while (r->has_chd()) for (Tri *c: r->ch)
                                                                       if (c && c->contains(p)) { r = c; break; }
                                                                      array<Tri*, 3> t; /* split into 3 triangles */
6.26 3D Projection [68f350]
                                                                      F3 t[i] = new (it++) Tri(r->p[i], r->p[R(i)], p);
using P3F = valarray<llf>;
                                                                      F3 link(E(t[i], 0), E(t[R(i)], 1));
F3 link(E(t[i], 2), r->e[L(i)]);
P3F toP3F(P3 p) { return {p.x, p.y, p.z}; }
llf dot(P3F a, P3F b) {
                                                                      r->ch = t:
return a[0]*b[0]+a[1]*b[1]+a[2]*b[2];
                                                                      F3 flip(t[i], 2);
P3F housev(P3 A, P3 B, int s) {
  const llf a = abs(A), b = abs(B);
  return toP3F(A) / a + s * toP3F(B) / b;
                                                                     auto build(const vector<P> &p) {
                                                                      it = pool; int n = (int)p.size();
                                                                      vector<int> ord(n); iota(all(ord), 0);
                                                                      shuffle(all(ord), mt19937(114514));
P project(P3 p, P3 q) {
                                                                      root = new (it++) Tri(n, n + 1, n + 2);
P3 o(0, 0, 1);
                                                                      copy_n(p.data(), n, v); v[n++] = P(-C, -C);
v[n++] = P(C * 2, -C); v[n++] = P(-C, C * 2);
P3F u = housev(q, o, q.z > 0 ? 1 : -1);
 auto pf = toP3F(p);
                                                                      for (int i : ord) add_point(i);
 auto np = pf - 2 * u * dot(u, pf) / dot(u, u);
                                                                      vector<array<int, 3>> res;
 return P(np[0], np[1]);
                                                                      for (Tri *now = pool; now != it; now++)
} // project p onto the plane q^Tx = 0
                                                                       if (!now->has_chd()) res.push_back(now->p);
6.27 3D Skew Line Nearest Point
                                                                      return res:
• L_1: \mathbf{v}_1 = \mathbf{p}_1 + t_1 \mathbf{d}_1, L_2: \mathbf{v}_2 = \mathbf{p}_2 + t_2 \mathbf{d}_2
• n = d_1 \times d_2
                                                                     6.29 Build Voronoi [94f000]
• \boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}, \boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}
                                                                     void build_voronoi_cells(auto &&p, auto &&res) {
m{\cdot} \ \ m{c}_1 = m{p}_1 + rac{(m{p}_2 - m{p}_1) \cdot m{n}_2}{m{d}_1 \cdot m{n}_2} m{d}_1, m{c}_2 = m{p}_2 + rac{(m{p}_1 - m{p}_2) \cdot m{n}_1}{m{d}_2 \cdot m{n}_1} m{d}_2
                                                                      vector<vector<int>> adj(p.size());
6.28 Delaunay [3a4ff1] - 1aee24/19ec42
                                                                      for (auto f: res) F3 {
/* please ensure input points are unique *,
                                                                       int a = f[i], b = f[R(i)];
                                                                       if (a >= p.size() || b >= p.size()) continue;
/* A triangulation such that no points will strictly
inside circumcircle of any triangle. C should be big
                                                                       adj[a].emplace_back(b);
enough s.t. the initial triangle contains all points */
                                                                      // use `adj` and `p` and HPI to build cells
for (size_t i = 0; i < p.size(); i++) {</pre>
#define L(i) ((i)==0 ? 2 : (i)-1)
#define R(i) ((i)==2 ? 0 : (i)+1)
#define F3 for (int i = 0; i < 3; i++)
                                                                       vector<Line> ls = frame; // the frame
bool is_inf(P z) { return RE(z) <= -C || RE(z) >= C; }
                                                                       for (int j : adj[i]) {
                                                                        P = p[i] + p[j], d = rot90(p[j] - p[i]);
bool in_cc(const array<P,3> &p, P q) {
 i128 inf_det = 0, det = 0, inf_N, N;
                                                                        assert (norm(d) != 0);
                                                                        ls.emplace_back(m, m + d); // doubled coordinate
 F3 {
                                                                       } // HPI(ls)
  if (is_inf(p[i]) && is_inf(q)) continue;
 else if (is_inf(p[i])) inf_N = 1, N = -norm(q);
else if (is_inf(q)) inf_N = -1, N = norm(p[i]);
else inf_N = 0, N = norm(p[i]) - norm(q);
                                                                     6.30 kd Tree (Nearest Point)* [f733e5]
  lld D = cross(p[R(i)] - q, p[L(i)] - q);
                                                                     struct KDTree {
  inf_det += inf_N * D; det += N * D;
                                                                      struct Node {
                                                                       int x, y, x1, y1, x2, y2, id, f; Node *L, *R;
 return inf_det != 0 ? inf_det > 0 : det > 0;
                                                                      } tree[maxn], *root;
                                                                      lld dis2(int x1, int y1, int x2, int y2) {
                                                                       lld dx = x1 - x2, dy = y1 - y2;
P v[maxn];
struct Tri;
                                                                       return dx * dx + dy * dy;
struct E {
                                                                      static bool cmpx(Node& a, Node& b) { return a.x<b.x; }</pre>
 Tri *t; int side;
 E(Tri *t_=0, int side_=0) : t(t_), side(side_) {}
                                                                      static bool cmpy(Node& a, Node& b) { return a.y<b.y; }</pre>
                                                                      void init(vector<pair<int,int>> &ip) {
                                                                       for (int i = 0; i < ssize(ip); i++)
tie(tree[i].x, tree[i].y) = ip[i], tree[i].id = i;</pre>
struct Tri {
 array<int,3> p; array<Tri*,3> ch; array<E,3> e;
 Tri(int a=0, int b=0, int c=0) : p{a, b, c}, ch{} {}
                                                                       root = build(0, (int)ip.size()-1, 0);
 bool has_chd() const { return ch[0] != nullptr; }
 bool contains(int q) const {
                                                                      Node* build(int L, int R, int d) {
 F3 if (ori(v[p[i]], v[p[R(i)]], v[q]) < 0)
                                                                       if (L>R) return nullptr;
   return false;
                                                                       int M = (L+R)/2;
  return true;
                                                                       nth_element(tree+L,tree+M,tree+R+1,d%2?cmpy:cmpx);
                                                                       Node &o = tree[M]; o.f = d \% 2;
bool check(int q) const {
                                                                       o.x1 = o.x2 = o.x; o.y1 = o.y2 = o.y;
return in_cc({v[p[0]], v[p[1]], v[p[2]]}, v[q]); }
pool[maxn * 10], *it, *root;
                                                                       o.L = build(L, M-1, d+1); o.R = build(M+1, R, d+1);
                                                                       for (Node *s: {o.L, o.R}) if (s) {
/* SPLIT_HASH_HERE */
                                                                        o.x1 = min(o.x1, s->x1); o.x2 = max(o.x2, s->x2);
void link(const E &a, const E &b) {
                                                                        o.y1 = min(o.y1, s->y1); o.y2 = max(o.y2, s->y2);
 if (a.t) a.t->e[a.side] = b;
if (b.t) b.t->e[b.side] = a;
                                                                       return tree+M;
                                                                      bool touch(int x, int y, lld d2, Node *r){
void flip(Tri *A, int a) {
auto [B, b] = A->e[a]; /* flip edge between A,B */
                                                                       lld d = (lld) \operatorname{sqrt}(d2) + 1;
 if (!B || !A->check(B->p[b])) return;
                                                                       return x \ge r - x_1 - d \& x \le r - x_2 + d \& x
 Tri *X = new (it++) Tri(A->p[R(a)], B->p[b], A->p[a]);
                                                                               y >= r->y1 - d \&\& y <= r->y2 + d;
Tri *Y = new (it++) Tri(B->p[R(b)], A->p[a], B->p[b]);
```

```
using P = pair<lld, int>;
 void dfs(int x, int y, P &mn, Node *r) {
  if (!r || !touch(x, y, mn.first, r)) return;
  mn = min(mn, P(dis2(r->x, r->y, x, y), r->id));
if (r->f == 1 ? y < r->y : x < r->x)
   dfs(x, y, mn, r\rightarrow L), dfs(x, y, mn, r\rightarrow R);
  else
   dfs(x, y, mn, r\rightarrow R), dfs(x, y, mn, r\rightarrow L);
 int query(int x, int y) {
  P mn(INF, -1); dfs(x, y, mn, root);
  return mn.second;
 }
} tree;
6.31 kd Closest Pair (3D ver.)* [84d9eb]
llf solve(vector<P> v) {
 shuffle(v.begin(), v.end(), mt19937());
 unordered_map<lld, unordered_map<lld,</pre>
  unordered_map<lld, int>>> m;
 llf d = dis(v[0], v[1]);
 auto Idx = [\&d] (Ilf x) \rightarrow Ild {
  return round(x * 2 / d) + 0.1; };
 auto rebuild_m = [&m, &v, &Idx](int k) {
  m.clear();
  for (int i = 0; i < k; ++i)</pre>
   m[Idx(v[i].x)][Idx(v[i].y)]
    [Idx(v[i].z)] = i;
 }; rebuild_m(2);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
     kz = Idx(v[i].z); bool found = false;
  for (int dx = -2; dx \le 2; ++dx) {
   const lld nx = dx + kx;
   if (m.find(nx) == m.end()) continue;
   auto& mm = m[nx];
   for (int dy = -2; dy <= 2; ++dy) {
    const lld ny = dy + ky;
    if (mm.find(ny) == mm.end()) continue;
    auto& mmm = mm[ny];
    for (int dz = -2; dz <= 2; ++dz) {
     const lld nz = dz + kz;
     if (mmm.find(nz) == mmm.end()) continue;
     const int p = mmm[nz];
     if (dis(v[p], v[i]) < d) {</pre>
      d = dis(v[p], v[i]);
      found = true;
     }
  if (found) rebuild_m(i + 1);
  else m[kx][ky][kz] = i;
 return d;
6.32 Simulated Annealing* [4e0fe5]
llf anneal() {
 mt19937 rnd_engine(seed);
 uniform_real_distribution<llf> rnd(0, 1);
 const llf dT = 0.001;
 // Argument p
llf S_cur = calc(p), S_best = S_cur;
for (llf T = 2000; T > EPS; T -= dT) {
  // Modify p to p_prime
  const llf S_prime = calc(p_prime);
  const llf delta_c = S_prime - S_cur;
  llf prob = min((llf)1, exp(-delta_c / T));
  if (rnd(rnd_engine) <= prob)</pre>
   S_cur = S_prime, p = p_prime;
  if (S_prime < S_best) // find min</pre>
   S_best = S_prime, p_best = p_prime;
return S_best;
6.33 Triangle Centers* [adb146]
0 = ... // see min circle cover
G = (A + B + C) / 3;
H = G * 3 - 0 * 2; // orthogonal center
llf a = abs(B - C), b = abs(A - C), c = abs(A - B);
I = (a * A + b * B + c * C) / (a + b + c);
// FermatPoint: minimizes sum of distance
```

```
// I max. angle >= 120 deg then vertex
// otherwise, make eq. triangle AB'C, CA'B, BC'A
// line AA', BB', CC' intersects at P

Stringology
      Stringology
     Hash [ce7fad]
template <int P = 127, int Q = 1051762951>
class Hash {
 vector<int> h, p;
public:
 Hash(const auto &s) : h(s.size()+1), p(s.size()+1) {
  for (size_t i = 0; i < s.size(); ++i)</pre>
   h[i + 1] = add(mul(h[i], P), s[i]);
  generate(all(p), [x = 1, y = 1, this]() mutable {
   return y = x, x = mul(x, P), y; });
 int query(int l, int r) const { // 0-base [l, r)
  return sub(h[r], mul(h[l], p[r - l]));
};
7.2 Suffix Array [ald8fe] - 9603dl/eb7a2f
auto sais(const auto &s) {
 const int n = (int)s.size(), z = ranges::max(s) + 1;
 if (n == 1) return vector{0};
 vector<int> c(z); for (int x : s) ++c[x];
 partial_sum(all(c), begin(c));
 vector<int> sa(n); auto I = views::iota(0, n);
 vector<bool> t(n); t[n - 1] = true;
for (int i = n - 2; i >= 0; --i)
  t[i] = (s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 auto is_lms = views::filter([&t](int x) {
  return x && t[x] && !t[x - 1]; });
 auto induce = [&] {
  for (auto x = c; int y : sa)
   if (y--) if (!t[y]) sa[x[s[y] - 1]++] = y;
  for (auto x = c; int y : sa | views::reverse)
   if (y--) if (t[y]) sa[--x[s[y]]] = y;
 vector<int> lms, q(n); lms.reserve(n);
 for (auto x = c; int i : I | is_lms) {
  q[i] = int(lms.size());
  lms.push_back(sa[--x[s[i]]] = i);
 induce(); vector<int> ns(lms.size());
 for (int j = -1, nz = 0; int i : sa | is_lms) {
  if (j > = 0) {
   int len = min({n - i, n - j, lms[q[i] + 1] - i});
    ns[q[i]] = nz += lexicographical_compare(
      begin(s) + j, begin(s) + j + len,
begin(s) + i, begin(s) + i + len);
  j = i;
 }
 ranges::fill(sa, 0); auto nsa = sais(ns);
 for (auto x = c; int y : nsa | views::reverse)
  y = lms[y], sa[--x[s[y]]] = y;
 return induce(), sa;
// SPLIT_HASH_HERE sa[i]: sa[i]-th suffix is the
// i-th lexicographically smallest suffix.
// hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
struct Suffix {
 int n; vector<int> sa, hi, rev;
 Suffix(const auto &s) : n(int(s.size())),
  hi(n), rev(n) {
  vector<int> _s(n + 1); // _s[n] = 0;
copy(all(s), begin(_s)); // s shouldn't contain 0
  sa = sais(_s); sa.erase(sa.begin());
  for (int i = 0; i < n; ++i) rev[sa[i]] = i;
for (int i = 0, h = 0; i < n; ++i) {</pre>
   if (!rev[i]) { h = 0; continue; }
   for (int j = sa[rev[i] - 1]; i + h < n && j + h < n</pre>
   && s[i + h] == s[j + h];) ++h;
hi[rev[i]] = h ? h-- : 0;
 }
7.3 Suffix Array Tools* [8e08c8]
template <int LG = 20> struct SparseTableSA : Suffix {
 array<vector<int>, LG> mn;
 SparseTableSA(const auto &s) : Suffix(s), mn{hi} {
  for (int l = 0; l + 1 < LG; l++) { mn[l+1].resize(n);</pre>
```

```
for (int i = 0, len = 1 << l; i + len < n; i++)</pre>
    mn[l + 1][i] = min(mn[l][i], mn[l][i + len]);
                                                               return f;
  }
                                                              }
                                                              vector<int> search(const auto &s, const auto &t) {
 int lcp(int a, int b) {
                                                               // return 0-indexed occurrence of t in s
                                                               vector<int> f = kmp(t), r;
for (int i = 0, k = 0; i < (int)s.size(); ++i) {
  while (k > 0 && s[i] != t[k]) k = f[k - 1];
  if (a == b) return n - a;
  a = rev[a] + 1, b = rev[b] + 1;
 if (a > b) swap(a, b);
  const int lg = __lg(b - a);
                                                                k += (s[i] == t[k]);
  return min(mn[lg][a], mn[lg][b - (1 << lg)]);</pre>
                                                                if (k == (int)t.size())
                                                                 r.push_back(i - t.size() + 1), k = f[k - 1];
 } // equivalent to lca on the kruskal tree
 pair<int,int> get_range(int x, int len) { // WIP
  int a = rev[x] + 1, b = rev[x] + 1;
                                                               return r:
  for (int l = LG - 1; l >= 0; l--) {
   const int s = 1 << l;
                                                              7.6
                                                                   Z value [6a7fd0]
   if (a + s <= n && mn[l][a] >= len) a += s;
                                                              vector<int> Zalgo(const string &s) {
   if (b - s >= 0 && mn[l][b - s] >= len) b -= s;
                                                               vector<int> z(s.size(), s.size());
                                                               for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
  int j = clamp(r - i, 0, z[i - l]);</pre>
  return {b - 1, a};
 } // if offline, solve get_range with DSU
                                                                for (; i + j < z[0] \text{ and } s[i + j] == s[j]; ++j);
};
7.4
                                                                if (i + (z[i] = j) > r) r = i + z[l = i];
     Ex SAM* [58374b]
struct exSAM {
 int len[maxn * 2], link[maxn * 2]; // maxlen, suflink
 int next[maxn * 2][maxc], tot; // [0, tot), root = 0
                                                              7.7 Manacher [c938a9]
 int ord[maxn * 2]; // topo. order (sort by length)
                                                              vector<int> manacher(const string &S) {
 int cnt[maxn * 2]; // occurence
                                                               const int n = (int)S.size(), m = n * 2 + 1;
 int newnode() {
                                                               vector<int> z(m);
  fill_n(next[tot], maxc, 0);
                                                               string t = "."; for (char c: S) t += c, t += '.';
for (int i = 1, l = 0, r = 0; i < m; ++i) {
  return len[tot] = cnt[tot] = link[tot] = 0, tot++;
                                                                z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
 void init() { tot = 0, newnode(), link[0] = -1; }
                                                                while (i - z[i] >= 0 \&\& i + z[i] < m) {
 int insertSAM(int last, int c) {
                                                                 if (t[i - z[i]] == t[i + z[i]]) ++z[i];
  int cur = next[last][c];
                                                                 else break:
  len[cur] = len[last] + 1;
  int p = link[last];
                                                                if (i + z[i] > r) r = i + z[i], l = i;
  while (p != -1 && !next[p][c])
   next[p][c] = cur, p = link[p];
                                                               return z; // the palindrome lengths are z[i] - 1
  if (p == -1) return link[cur] = 0, cur;
  int q = next[p][c];
                                                              /* for (int i = 1; i + 1 < m; ++i) {
  if (len[p] + 1 == len[q]) return link[cur] = q, cur;
                                                                int l = (i - z[i] + 2) / 2, r = (i + z[i]) / 2;
  int clone = newnode();
                                                                if (l != r) // [l, r) is maximal palindrome
  for (int i = 0; i < maxc; ++i)</pre>
  next[clone][i] = len[next[q][i]] ? next[q][i] : 0;
                                                              7.8 Lyndon Factorization [d22cc9]
  len[clone] = len[p] + 1;
                                                              // partition s = w[0] + w[1] + ... + w[k-1],
  while (p != -1 && next[p][c] == q)
                                                              // w[0] >= w[1] >= ... >= w[k-1]
   next[p][c] = clone, p = link[p];
                                                              // each w[i] strictly smaller than all its suffix
  link[link[cur] = clone] = link[q];
                                                              void duval(const auto &s, auto &&report) {
  link[q] = clone;
                                                               for (int n = (int)s.size(), i = 0, j, k; i < n; ) {
  for (j = i + 1, k = i; j < n && s[k] <= s[j]; j++)</pre>
  return cur;
                                                                 k = (s[k] < s[j] ? i : k + 1);
 void insert(const string &s) {
                                                                // if (i < n / 2 && j >= n / 2) {
  int cur = 0:
                                                                // for min cyclic shift, call duval(s + s)
  for (char ch : s) {
                                                                // then here s.substr(i, n / 2) is min cyclic shift
   int &nxt = next[cur][int(ch - 'a')];
   if (!nxt) nxt = newnode();
                                                                for (; i <= k; i += j - k)
   cnt[cur = nxt] += 1;
                                                                 report(i, j - k); // s.substr(l, len)
  }
                                                              } // tested @ luogu 6114, 1368 & UVA 719
 void build() {
                                                              7.9 Main Lorentz* [615b8f]
  queue<int> q; q.push(0);
                                                              vector<pair<int, int>> rep[kN]; // 0-base [l, r]
  while (!q.empty()) {
   int cur = q.front(); q.pop();
                                                              void main_lorentz(const string &s, int sft = 0) {
   for (int i = 0; i < maxc; ++i)</pre>
                                                               const int n = s.size();
    if (next[cur][i]) q.push(insertSAM(cur, i));
                                                               if (n == 1) return;
                                                               const int nu = n / 2, nv = n - nu;
  }
                                                               const string u = s.substr(0, nu), v = s.substr(nu)
  vector<int> lc(tot);
                                                               ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend());
main_lorentz(u, sft), main_lorentz(v, sft + nu);
  for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
  partial_sum(all(lc), lc.begin());
  for (int i = 1; i < tot; ++i) ord[--lc[len[i]]] = i;</pre>
                                                               const auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u),
                                                                      z3 = Zalgo(ru + '#' + rv), z4 = Zalgo(v);
 void solve() {
                                                               auto get_z = [](const vector<int> &z, int i) {
  for (int i = tot - 2; i >= 0; --i)
                                                                return (0 <= i and i < (int)z.size()) ? z[i] : 0; };</pre>
                                                               auto add_rep = [&](bool left, int c, int l, int k1,
   cnt[link[ord[i]]] += cnt[ord[i]];
                                                                   int k2) {
                                                                const int L = max(1, l - k2), R = min(l - left, k1);
    KMP [3727f3]
                                                                if (L > R) return;
vector<int> kmp(const auto &s) {
                                                                if (left) rep[l].emplace_back(sft + c - R, sft + c -
vector<int> f(s.size());
                                                                  L);
 for (int i = 1, k = 0; i < (int)s.size(); ++i) {</pre>
                                                                else rep[l].emplace_back(sft + c - R - l + 1, sft + c
  while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
                                                                    - L - l + 1);
  f[i] = (k += (s[i] == s[k]));
                                                               };
```

```
for (int cntr = 0; cntr < n; cntr++) {</pre>
  int l, k1, k2;
  if (cntr < nu) {</pre>
   l = nu - cntr;
   k1 = get_z(z1, nu - cntr);
   k2 = get_z(z2, nv + 1 + cntr);
  } else {
   l = cntr - nu + 1;
   k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
   k2 = get_z(z4, (cntr - nu) + 1);
  if (k1 + k2 >= 1)
   add_rep(cntr < nu, cntr, l, k1, k2);</pre>
}
7.10 BWT* [a8287e]
```

```
void BWT(char *ori, char *res) {
  // make ori -> ori + ori then build suffix array
void iBWT(char *ori, char *res) {
 vector<int> v[SIGMA], a;
 const int len = strlen(ori); res[len] = 0;
 for (int i = 0; i < len; i++) v[ori[i] - 'a'].pb(i);</pre>
 for (int i = 0, ptr = 0; i < SIGMA; i++)</pre>
 for (int j : v[i]) a.pb(j), ori[ptr++] = 'a' + i;
for (int i = 0, ptr = 0; i < len; i++)</pre>
  res[i] = ori[a[ptr]], ptr = a[ptr];
```

## 7.11 Palindromic Tree\* [0673ee]

```
struct PalindromicTree {
struct node {
  int nxt[26], f, len; // num = depth of fail link
                 // = #pal_suffix of this node
  int cnt, num;
  node(int l = 0) : nxt{}, f(0), len(l), cnt(0), num(0) {}
};
vector<node> st; vector<char> s; int last, n;
void init() {
 st.clear(); s.clear(); last = 1; n = 0;
  st.push_back(0); st.push_back(-1);
  st[0].f = 1; s.push_back(-1);
int getFail(int x) {
 while (s[n - st[x].len - 1] != s[n]) x = st[x].f;
  return x;
void add(int c) {
    s.push_back(c -= 'a'); ++n;
  int cur = getFail(last);
  if (!st[cur].nxt[c]) {
   int now = st.size();
   st.push_back(st[cur].len + 2);
   st[now].f = st[getFail(st[cur].f)].nxt[c];
   st[cur].nxt[c] = now;
  st[now].num = st[st[now].f].num + 1;
  7
  last = st[cur].nxt[c]; ++st[last].cnt;
}
void dpcnt() { // cnt = #occurence in whole str
  for (int i = st.size() - 1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
int size() { return st.size() - 2; }
} pt; /* string s; cin >> s; pt.init();
for (int i = 0; i < SZ(s); i++)
int prvsz = pt.size(); pt.add(s[i]);
if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
  // pal @ [l,r]: s.substr(l, r-l+1)
} */
```

#### 8 Misc Theorems

#### Spherical Coordinate





$$\begin{split} r &= \sqrt{x^2 + y^2 + z^2} \\ \theta &= \mathsf{acos}(z/\sqrt{x^2 + y^2 + z^2}) \\ \phi &= \mathsf{atan2}(y,x) \end{split}$$

#### **Spherical Cap**

- · A portion of a sphere cut off by a plane.
- r: sphere radius, a: radius of the base of the cap, h: height of the cap,  $\theta$ :
- $\begin{aligned} & \text{Area} &= \pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta) (1 \cos \theta)^2/3. \\ & \text{Area} &= 2\pi r h = \pi (a^2 + h^2) = 2\pi r^2 (1 \cos \theta). \end{aligned}$

#### Sherman-Morrison formula

$$(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1+v^{\mathsf{T}}A^{-1}u}$$

Kirchhoff's Theorem Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} =$ 

- d(i),  $L_{ij} = -c$  where c is the number of edge (i, j) in G. • The number of undirected spanning in G is  $\det(\tilde{L}_{11})$ .
- The number of directed spanning tree rooted at r in G is  $\det(\tilde{L}_{rr})$ .

#### BEST Theorem

 $\#\{\text{Eulerian circuits}\} = \#\{\text{arborescences rooted at l}\} \cdot \prod_{v \in V} (\deg(v) - 1)!$ 

#### Random Walk on Graph

Let P be the transition matrix of a strongly connected directed graph,  $\sum_{i} P_{i,j} = 1$ . Let  $F_{i,j}$  be the expected time to reach j from i. Let  $g_i$  be the expected time from i to i, G = diag(g) and J be a matrix all of 1, i.e.  $J_{i,j} = 1$ . Then, F = J - G + PF

First solve G: let  $\pi P = \pi$  be a stationary distribution. Then  $\pi_i g_i = 1$ . The rank of I-P is n-1, so we first solve a special solution X such that (I-P)X=J-G and adjust X to F by  $F_{i,j}=X_{i,j}-X_{j,j}$ .

#### **Tutte Matrix**

For  $i < j, d_{ij} = x_{ij}$  (in practice, a random number) if  $(i,j) \in \mathit{E}$ , otherwise  $d_{ij}=0$ . For  $i\geq j, d_{ij}=-d_{ji}$ .  $\frac{\mathrm{rank}(D)}{2}$  is the maximum matching.

## Cayley's Formula

- · Given a degree sequence  $d_1, d_2, \ldots, d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(n-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=$  $kn^{n-k-1}$ .

#### Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and  $\sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k)$  holds for all  $1 \le k \le n$ .

### Havel-Hakimi algorithm

Find the vertex who has greatest degree unused, connect it with other greatest vertex.

#### Gale-Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \geq \cdots \geq a_n$  and  $b_1, \ldots, b_n$  is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i, k)$  holds for every  $1 \leq k \leq n$ .

#### Fulkerson-Chen-Anstee theorem

A sequence  $(a_1,b_1),\ldots,(a_n,b_n)$  of nonnegative integer pairs with  $a_1\geq \cdots \geq a_n$  is digraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \leq a_i$  $\sum_{i=1}^k \min(b_i, k-1) + \sum_{i=k+1}^n \min(b_i, k)$  holds for every  $1 \le k \le n$ . **Euler's planar graph formula** 

V - E + F = C + 1.  $E \le 3V - 6$  (when  $V \ge 3$ )

#### Pick's theorem

For simple polygon, when points are all integer, we have  ${\cal A}$  $\#\{\text{lattice points in the interior}\} + \frac{1}{2}\#\{\text{lattice points on the boundary}\} - 1$ 

### Matroid

- $\begin{array}{l} \boldsymbol{\cdot} \;\; B\subseteq A \wedge A \in \mathcal{I} \Rightarrow B \in \mathcal{I}. \\ \boldsymbol{\cdot} \;\; \text{If } A,B \in \mathcal{I} \text{ and } |A| > |B| \text{, then } \exists x \in A \setminus B, B \cup \{x\} \in \mathcal{I}. \end{array}$

 $A \in I$  iff linear indep. Linear matroid  $I={\sf forests}$  of undirected graph Graphic matroid Colorful matroid (EX) Each color c has an upper bound  $R_c$ Transversal matroid  $A \in I$  iff  $\exists$  matching M whose right part is A.  $A \in I$  iff G is connected after removing edges A. Bond matroid  $A \in I^*$  iff there is a basis  $\subseteq E \setminus A$   $A \in I'$  iff  $A \in I \land |A| \le k$ **Dual matroid** Truncated matroid

## **Matroid Intersection**

Given matroids  $M_1=(G,I_1), M_2=(G,I_2)$ , find maximum  $S\in I_1\cap I_2$ . For each iteration, build the directed graph and find a shortest path from s to t.

•  $s \rightarrow x : S \sqcup \{x\} \in I_1$ 

 $\begin{array}{l} \cdot \ x \to t : S \sqcup \{x\} \in I_2 \\ \cdot \ y \to x : S \setminus \{y\} \sqcup \{x\} \in I_1 \ (y \ \text{is in the unique circuit of} \ S \sqcup \{x\} \} \\ \cdot \ x \to y : S \setminus \{y\} \sqcup \{x\} \in I_2 \ (y \ \text{is in the unique circuit of} \ S \sqcup \{x\} \} \end{array}$ 

Alternate the path, and |S| will increase by 1. In each iteration, |E| = O(RN), where  $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)), N=|G|$ . For weighted case, assign weight -w(x) and w(x) to  $x\in S$  and  $x\notin S$ , resp. Find the shortest path by Bellman-Ford. The maximum iteration of Bellman-Ford is 2R+1.

#### Dual of LP

Primal	Dual
Maximize $c^{T}x$ s.t. $Ax \leq b, x \geq 0$	Minimize $b^{T}y$ s.t. $A^{T}y \geq c$ , $y \geq 0$
Maximize $c^{T}x$ s.t. $Ax \leq b$	Minimize $b^{T}y$ s.t. $A^{T}y = c$ , $y \geq 0$
Maximize $c^{T}x$ s.t. $Ax = b, x \geq 0$	Minimize $b^{T}y$ s.t. $A^{T}y \geq c$

## **Dual of Min Cost b-Flow**

- Capacity  $c_{uv}$ , Flow  $f_{uv}$ , Cost  $w_{uv}$ , Required Flow difference for vertex  $b_u$ .
- If all  $w_{uv}$  are integers, then optimal solution can happen when all  $p_u$  are integers.

$$\begin{split} \min \sum_{uv} w_{uv} f_{uv} \text{ s.t. } -f_{uv} &\geq -c_{uv}, \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_u \\ \Leftrightarrow \min \sum_{u} b_u p_u + \sum_{uv} c_{uv} \max(0, p_v - p_u - w_{uv}) \text{ s.t. } p_u &\geq 0 \end{split}$$

```
Minimax Theorem
```

Let  $f: X \times Y \to \mathbb{R}$  be continuous where  $X \subseteq \mathbb{R}^n, Y \subseteq \mathbb{R}^m$  are compact and convex. If  $f(\cdot,y): X \to \mathbb{R}$  is concave for fixed y, and  $f(x,\cdot): Y \to \mathbb{R}$  is convex for fixed x, then  $\max_{x \in X} \min_{y \in Y} f(x,y) = \min_{y \in Y} \max_{x \in X} f(x,y)$ , e.g.  $f(x,y) = x^\intercal Ay$  for zero-sum matrix game.

#### Parallel Axis Theorem

12: end while

The second moment of area is  $I_z=\iint x^2+y^2\mathrm{d}A.$   $I_{z'}=I_z+Ad^2$  where d is the distance between two parallel axis z,z'.

## 8.2 Stable Marriage

```
l: Initialize m \in M and w \in W to free

2: while \exists free man m who has a woman w to propose to do

3: w \leftarrow first woman on m's list to whom m has not yet proposed

4: if \exists some pair (m', w) then

5: if w prefers m to m' then

6: m' \leftarrow free

7: (m, w) \leftarrow engaged

8: end if

9: else

10: (m, w) \leftarrow engaged

11: end if
```

```
8.3 Weight Matroid Intersection* [d00ee8]
struct Matroid {
 Matroid(bitset<N>); // init from an independent set
 bool can_add(int); // check if break independence
 Matroid remove(int); // removing from the set
auto matroid_intersection(const vector<int> &w) {
 const int n = (int)w.size(); bitset<N> S;
 for (int sz = 1; sz <= n; sz++) {</pre>
 Matroid M1(S), M2(S); vector<vector<pii>>> e(n + 2);
  for (int j = 0; j < n; j++) if (!S[j]) {</pre>
   if (M1.can_add(j)) e[n].eb(j, -w[j]);
   if (M2.can_add(j)) e[j].eb(n + 1, 0);
  for (int i = 0; i < n; i++) if (S[i]) {</pre>
   Matroid T1 = M1.remove(i), T2 = M2.remove(i);
   for (int j = 0; j < n; j++) if (!S[j]) {
  if (T1.can_add(j)) e[i].eb(j, -w[j]);
  if (T2.can_add(j)) e[i].eb(j, -w[j]);</pre>
    if (T2.can_add(j)) e[j].eb(i, w[i]);
  } // maybe implicit build graph for more speed
  vector<pii> d(n + 2, {INF, 0}); d[n] = {0, 0};
  vector<int> prv(n + 2, -1);
  // change to SPFA for more speed, if necessary
  for (int upd = 1; upd--; )
   for (int u = 0; u < n + 2; u++)
    for (auto [v, c] : e[u]) {
     pii x(d[u].first + c, d[u].second + 1);
     if (x < d[v]) d[v] = x, prv[v] = u, upd = 1;
  if (d[n + 1].first >= INF) break;
  for (int x = prv[n+1]; x!=n; x = prv[x]) S.flip(x);
  // S is the max-weighted independent set w/ size sz
 return S;
} // from Nacl
8.4 Bitset LCS [4155ab]
cin >> n >> m;
for (int i = 1, x; i <= n; ++i)</pre>
cin >> x, p[x].set(i);
for (int i = 1, x; i <= m; ++i) {
  cin >> x, (g = f) |= p[x];
 f.shiftLeftByOne(), f.set(0);
 ((f = g - f) ^= g) \&= g;
cout << f.count() << '\n';</pre>
      Prefix Substring LCS [7d8faf]
8.5
void all_lcs(string S, string T) { // 0-base
 vector<size_t> h(T.size()); iota(all(h), 1);
 for (size_t a = 0; a < S.size(); ++a) {</pre>
  for (size_t c = 0, v = 0; c < T.size(); ++c)</pre>
   if (S[a] == T[c] || h[c] < v) swap(h[c], v);</pre>
  // here, LCS(s[0, a], t[b, c]) =
  // c - b + 1 - sum([h[i] > b] | i \le c)
   // test @ yosupo judge

Convex 1D/1D DP* [938911]
struct S { int i, l, r; };
auto solve(int n, int k, auto &w) {
vector<int64_t> dp(n + 1); dp[0] = 0;
auto f = [&](int l, int r) -> int64_t {
```

if  $(r - \bar{l} > k)$  return -INF;

```
return dp[l] + w(l + 1, r);
 deque<S> dq; dq.emplace_back(0, 1, n);
 for (int i = 1; i <= n; ++i) {</pre>
  dp[i] = f(dq.front().i, i);
  while (!dq.empty() && dq.front().r <= i)</pre>
   dq.pop_front();
  dq.front().l = i + 1;
  while (!dq.empty() &&
    f(i, dq.back().l) >= f(dq.back().i, dq.back().l))
   dq.pop_back();
  int p = i + 1;
  if (!dq.empty()) {
   auto [j, l, r] = dq.back();
   for (int s = 1 << 20; s; s >>= 1)
    if (l+s <= n && f(i, l+s) < f(j, l+s)) l += s;</pre>
   dq.back().r = l; p = l + 1;
  if (p <= n) dq.emplace_back(i, p, n);</pre>
 return dp;
} // test @ tioj 烏龜疊疊樂
8.7 ConvexHull Optimization [b4318e]
struct L {
 mutable lld a, b, p;
 bool operator<(const L &r) const {</pre>
  return a < r.a; /* here */ }
 bool operator<(lld x) const { return p < x; }</pre>
lld Div(lld a, lld b) {
 return a / b - ((a ^ b) < 0 && a % b); }
struct DynamicHull : multiset<L, less<>>> {
 static const lld kInf = 1e18;
 bool Isect(iterator x, iterator y) {
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a)
   x->p = x->b > y->b ? kInf : -kInf; /* here */
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(lld a, lld b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
  while ((y = x) != begin() && (--x)->p >= y->p)
   Isect(x, erase(y));
 lld Query(lld x) { // default chmax
  auto l = *lower_bound(x); // to chmin:
  return l.a * x + l.b; // modify the 2 "<>"
};
8.8 Min Plus Convolution [464dcd]
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector<int> min_plus_convolution(auto &a, auto &b) {
 const int n = (int)a.size(), m = (int)b.size();
 vector<int> c(n + m - 1, numeric_limits<int>::max());
 auto dc = [&](auto Y, int l, int r, int jl, int jr) {
  if (l > r) return;
  int mid = (l + r) / 2, from = -1, &best = c[mid];
  for (int j = jl; j <= jr; j++)</pre>
   if (int i = mid - j; i >= 0 && i < n)
    if (best > a[i]+b[j]) best = a[i]+b[j], from = j;
  Y(Y, l, mid-1, jl, from); Y(Y, mid+1, r, from, jr);
 };
 return dc(dc, 0, n-1+m-1, 0, m-1), c;
8.9 SMAWK [f37761]
// For all 2x2 submatrix:
  If M[1][0] < M[1][1], M[0][0] < M[0][1]
// If M[1][0] == M[1][1], M[0][0] <= M[0][1]
// M[i][ans_i] is the best value in the i-th row
VI smawk(int N, int M, auto &&select) {
 auto dc = [&](auto self, const VI &r, const VI &c) {
  if (r.empty()) return VI{};
  const int n = (int)r.size(); VI ans(n), nr, nc;
  for (int i : c) {
   while (!nc.empty() &&
     select(r[nc.size() - 1], nc.back(), i))
    nc.pop_back();
   if (int(nc.size()) < n) nc.push_back(i);</pre>
```

```
if (k \& 1) p.x = -p.x;
  for (int i = 1; i < n; i += 2) nr.push_back(r[i]);</pre>
                                                               else swap(p.x, p.y);
  const auto na = self(self, nr, nc);
                                                             }
  for (int i = 1; i < n; i += 2) ans[i] = na[i >> 1];
                                                             return edges; // [{w, i, j}, ...]
 for (int i = 0, j = 0; i < n; i += 2) {
                                                            } // test @ yosupo judge
  ans[i] = nc[j];
   const int end = i + 1 == n ? nc.back() : ans[i + 1];
   while (nc[j] != end)
    if (select(r[i], ans[i], nc[++j])) ans[i] = nc[j];
  return ans;
VI R(N), C(M); iota(all(R), 0), iota(all(C), 0);
                                                            8.14 Binary Search On Fraction [ff3abd]
return dc(dc, R, C);
                                                            struct Q {
                                                             lld p, q; // p / q
bool min_plus_conv_select(int r, int u, int v) {
                                                             Q go(Q b, lld d) { return {p + b.p*d, q + b.q*d}; }
auto f = [](int i, int j) {
   if (0 <= i - j && i - j < n) return b[j] + a[i - j];</pre>
                                                            // returns smallest p/q in [lo, hi] such that
  return 2100000000 + (i - j);
                                                            // pred(p/q) is true, and 0 <= p,q <= N
                                                            Q frac_bs(lld N, auto &&pred) {
return f(r, u) > f(r, v);
                                                             Q lo{0, 1}, hi{1, 0};
\} // if f(r, v) is better than f(r, u), return true
                                                             if (pred(lo)) return lo;
8.10 De-Bruijn [aa7700]
                                                             assert(pred(hi));
vector<int> de_bruijn(int k, int n) {
                                                             bool dir = 1, L = 1, H = 1;
// return cyclic string of len k^n s.t. every string
                                                              for (; L || H; dir = !dir) {
// of len n using k char appears as a substring.
                                                              lld len = 0, step = 1;
vector<int> aux(n + 1), res;
                                                              for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
 auto db = [&](auto self, int t, int p) -> void {
                                                               if (Q mid = hi.go(lo, len + step);
 if (t <= n)
                                                                 mid.p > N || mid.q > N || dir ^ pred(mid))
   for (int i = aux[t - p]; i < k; ++i, p = t)</pre>
 aux[t] = i, self(self, t + 1, p);
else if (n % p == 0) for (int i = 1; i <= p; ++i)</pre>
                                                               else len += step;
                                                              swap(lo, hi = hi.go(lo, len));
  res.push_back(aux[i]);
                                                              (dir ? L : H) = !!len;
};
return db(db, 1, 1), res;
                                                             return dir ? hi : lo;
8.11 Josephus Problem [7f9ceb]
lld f(lld n, lld m, lld k) { // n 人每次隔 m-1 個殺
lld s = (m - 1) % (n - k); // O(k)
for (lld i = n - k + 1; i <= n; i++) s = (s + m) % i;
return s;
lld kth(lld n, lld m, i128 k) { // died at kth
                                                            8.15 Cartesian Tree [2ed09d]
if (m == 1) return k;
                           // O(m log(n))
for (k = k*m+m-1; k >= n; k = k-n + (k-n)/(m-1));
                                                            auto CartesianTree(const auto &a) {
return k:
                                                             const int n = (int)a.size(); vector<int> pa(n+1, -1);
                                                             for (int i = 1; i < n; i++) {
} // k and result are 0-based, test @ CF 101955
8.12 N Queens Problem* [31f83e]
                                                              int &p = pa[i] = i - 1, l = n;
void solve(\overline{VI} &ret, int n) { // no sol when n=2,3
                                                              while (p != -1 && a[i] < a[p])
if (n % 6 == 2) {
                                                               tie(l, pa[l], p, pa[p]) = tuple(p, p, pa[p], i);
 for (int i = 2; i <= n; i += 2) ret.push_back(i);</pre>
                                                             return pa.pop_back(), pa;
 ret.push_back(3); ret.push_back(1);
  for (int i = 7; i <= n; i += 2) ret.push_back(i);</pre>
                                                            } // root is minimum
  ret.push_back(5);
} else if (n % 6 == 3) {
 for (int i = 4; i <= n; i += 2) ret.push_back(i);</pre>
 ret.push_back(2);
 for (int i = 5; i <= n; i += 2) ret.push_back(i);</pre>
 ret.push_back(1); ret.push_back(3);
} else {
 for (int i = 2; i <= n; i += 2) ret.push_back(i);</pre>
                                                            8.16 Nim Product [4ac1ce]
  for (int i = 1; i <= n; i += 2) ret.push_back(i);</pre>
                                                            #define rep(i, r) for (int i = 0; i < r; i++)
                                                            struct NimProd {
                                                             llu bit_prod[64][64]{}, prod[8][8][256][256]{};
       Manhattan MST [1008bc]
                                                             NimProd() {
vector<array<int, 3>> manhattanMST(vector<P> ps) {
                                                              rep(i, 64) rep(j, 64) if (i & j) {
vector<int> id(ps.size()); iota(all(id), 0);
                                                                int a = lowbit(i & j);
vector<array<int, 3>> edges;
                                                                bit_prod[i][j] = bit_prod[i ^ a][j] ^
for (int k = 0; k < 4; k++) {
                                                                bit_prod[(i ^ a) | (a-1)][(j ^ a) | (i & (a-1))];
 sort(all(id), [&](int i, int j) {
                                                              } else bit_prod[i][j] = 1ULL << (i | j);</pre>
  return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y; });</pre>
                                                              rep(e, 8) rep(f, 8) rep(x, 256) rep(y, 256)
 map<int, int> sweep;
                                                               rep(i, 8) if (x >> i & 1) rep(j, 8) if (y >> j & 1)
  for (int i : id) {
                                                                prod[e][f][x][y] ^= bit_prod[e * 8 + i][f * 8 + j];
   for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++))
                                                             llu operator()(llu a, llu b) const {
    if (P d = ps[i] - ps[it->second]; d.y > d.x) break;
                                                              llu r = 0;
    else edges.push_back({d.y + d.x, i, it->second});
                                                               rep(e, 8) rep(f, 8)
                                                               r ^= prod[e][f][a >> (e*8) & 255][b >> (f*8) & 255];
   sweep[-ps[i].y] = i;
                                                              return r;
                                                             }
  for (P &p : ps)
                                                            };
```

# 8.17 Grid