Contents

```
1 Basic
                              5.26 Adaptive Simpson . . 15
1.1
    vimrc .....
                              5.27 Golden Ratio Search.
                                                      15
    Debug Macro.....
1.2
                             6 Geometry
                                                      15
1.3
     Increase Stack .....
                                  Basic Geometry....
                                                      15
                              6.1
    Pragma Optimization
1.4
                          1
                                  2D Convex Hull .....
                                                       15
1.5
    IO Optimization....
                              6.3
                                  2D Farthest Pair ....
                                                       16
    SVG Writer .....
1.6
                                  MinMax Enclosing
2 Data Structure
                                  16
2.1
    Dark Magic . . . . . . . . . .
                          2
                                  Minkowski Sum....
                                                      16
                                  Segment Intersection 16
    Link-Cut Tree . . . . . .
2.2
                                  Half Plane Intersec-
                          2
2.3
    LiChao Segment Tree
                              6.7
                                                       16
                                  tion .....
2.4 Treap .....
                                  SegmentDist
                              6.8
2.5 Linear Basis ......
                          3
                                                      16
                                   (Sausage) . . . . . . . . . .
     Binary Search On
2.6
                                  Rotating Sweep Line
                                                       16
     3
                              6.10 Polygon Cut . . . . . . . .
                          3
3 Graph
                                  Point In Simple
    2-SAT (SCC).....
                                  Polygon............ 17
3.2
    BCC . . . . . . . . . .
                              6.12 Point In Hull (Fast) .. 17
    Round Square Tree..
3.3
                              6.13 Point In Polygon
3.4
    Edge TCC .....
                                  (Fast) . . . . . . . . . 17
35
    DMST.....
                              6.14 Tangent of Points
    Dominator Tree . . . . .
3.6
                                  To Hull .....
                                                      17
3.7
    Edge Coloring .....
                              6.15 Circle Class & Inter-
     Centroid Decom-
                                   section .....
     position.....
                              6.16 Circle Common
Tangent .....
                                                      18
3.9
    Lowbit Decompo-
                              6.17 Line-Circle Inter-
     sition .....
                                                      18
3.10 Virtual Tree.....
                                  section .....
                              6.18 Poly-Circle Inter-
3.11 Tree Hashing ......
                                  section .....
3.12 Mo's Algorithm on
                              6.19 Minimum Covering
6
                                  18
                              6.20 Circle Union .....
                                                       18
3.14 MaximalClique .....
                              6.21 Polygon Union..... 18
3.15 MaximumClique ....
                              6.22 3D Point .....
                                                       19
                              6.23 3D Convex Hull .....
3.16 Minimum Cycle . . . . .
                                                       19
                          7
                              6.24 3D Projection . . . . . .
                                                      19
                          7
4 Flow & Matching
                              6.25 3D Skew Line Near-
                                   est Point .....
4.1 HopcroftKarp......
                              6.26 Delaunay .....
                                                       19
     Kuhn Munkres.....
                              6.27 Build Voronoi . . . . . . 6.28 kd Tree (Nearest
                                                      20
4.3
     Flow Models.....
     Dinic .....
                                  Point) .....
    HLPP .....
4.5
                              6.29 kd Closest Pair (3D
     Global Min-Cut .....
 4.6
                                                      20
4.7
     GomoryHu Tree . . . . .
                                  Minimum Cost Max
Flow.....
                              6.30 Simulated Annealing
 4.8
                          9
                              6.31 Triangle Centers . . . .
                                                      20
4.9 Dijkstra Cost Flow . . .
                                                      21
                             7 Stringology
4.10 Minimum Cost Cir-
                                  Hash .....
              Graph
                                  Suffix Array . . . . . . . . . . . .
                                                       21
     Matching .....
                         10
                                  Ex SAM .....
4.12 Weighted Matching .
                         10
                                  11
5 Math
                              7.6
     Common Bounds ...
                              7.7
                                  Lyndon Factorization 22
5.2
    Equations . . . . . . . . . . . .
                                  Main Lorentz .....
5.3
    Extended FloorSum.
                                                      22
                                  BWT.....
5.4
    Integer Division . . . .
                              7.10 Palindromic Tree.... 22
    FloorSum .....
5.5
                             8 Misc
                                                      22
     ModMin .....
5.6
                                  Theorems . . . . . . . . . 22
5.7
     Floor Monoid Product
                         12
                                  Weight Matroid In-
                              8.2
5.8
    ax+by=gcd .....
                                                      23
                                  tersection .....
5.9
     Chinese Remainder .
                                                      23
                              8.3
                                  Stable Marriage . . . . .
5.10 DiscreteLog ......
                              8.4
                                  Bitset LCS .....
                                                      23
5.11
     Quadratic Residue ...
                              8.5 Prefix Substring LCS.
                                                      23
5.12 FWT.....
5.13 Packed FFT .....
                              8.6 Convex 1D/1D DP . . . .
                                                      23
                                  ConvexHull Opti-
5.14 CRT for arbitrary mod 12
                                  mization .....
5.15 NTT/FFT .....
                         13
                                  Min Plus Convolution 24
5.16 Formal Power Series
                              8.9 De-Bruijn ..... 24
5.17 Partition Number ...
5.18 Pi Count .....
                              8.10 Josephus Problem .. 24
                              8.11 N Queens Problem . .
5.19 Miller Rabin ......
                              8.12 Tree Knapsack ..... 24
5.20 Pollard Rho ......
5.21 Berlekamp Massey . .
                              8.13 Manhattan MST..... 24
5.22 Gauss Elimination... 14
                              8.14 Binary Search On
5.23 Charateristic Poly-
                                  Fraction.....
     Barrett Reduction . . .
8.16 Montgomery ......
5.25 Simplex Construction 15
```

```
Basic
1.1
     vimrc
se is nu ru et tgc sc hls cin cino+=j1 sw=2 sts=2 bs=2
    mouse=a "encoding=utf-8 ls=2
syn on | colo desert | filetype indent on
inoremap {<CR> {<CR>}<ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -g -std=gnu++20 -</pre>
    DCKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
    Wconversion -fsanitize=address,undefined,float-
    divide-by-zero, float-cast-overflow && echo success<
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -g -std=gnu</pre>
    ++20 && echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space
    :]' \| md5sum \| cut -c-6
let c_no_curly_error=1
" setxkbmap -option caps:ctrl_modifier
1.2 Debug Macro [a45c59]
#define all(x) begin(x), end(x)
#ifdef CKISEKI
#include <experimental/iterator>
#define safe cerr<<__PRETTY_FUNCTION__<<" line "<</pre>
     __LINE__<<" safe\n"
#define debug(a...) debug_(#a, a)
#define orange(a...) orange_(#a, a)
void debug_(auto s, auto ...a) {
  cerr << "\e[1;32m(" << s << ") = (";</pre>
 int f = 0;
 (..., (cerr << (f++ ? ", " : "") << a));
 cerr << ")\e[0m\n";</pre>
void orange_(auto s, auto L, auto R) {
  cerr << "\e[1;33m[ " << s << " ] = [ ";</pre>
 using namespace experimental;
 copy(L, R, make_ostream_joiner(cerr, ", "));
cerr << " ]\e[0m\n";</pre>
#else
#define safe ((void)0)
#define debug(...) safe
#define orange(...) safe
#endif
1.3 Increase Stack [b6856c]
const int size = 256 << 20;</pre>
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
       _("movq %0, %%rsp\n"::"r"(bak));
  asm
1.4 Pragma Optimization [6006f6]
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse, sse2, sse3, ssse3, sse4")
#pragma GCC target("popent,abm,mmx,avx,arch=skylake")
  _builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
1.5 IO Optimization [c9494b]
static inline int gc() {
 constexpr int B = 1<<20; static char buf[B], *p, *q;</pre>
 if (p == q) q = (p = buf) + fread(buf, 1, B, stdin);
 return q == buf ? EOF : *p++;
1.6 SVG Writer [57436c]
class SVG {
 void p(string_view s) { o << s; }</pre>
 void p(string_view s, auto v, auto... vs) {
  auto i = s.find('$');
  o << s.substr(0, i) << v, p(s.substr(i + 1), vs...);
 ofstream o; string c = "red";
public:
 SVG(auto f,auto x1,auto y1,auto x2,auto y2) : o(f) {
  p("<svg xmlns='http://www.w3.org/2000/svg'</pre>
   "viewBox='$ $ $'>\n"
   "<style>*{stroke-width:0.5%;}</style>\n",
 x1, -y2, x2 - x1, y2 - y1); }
~SVG() { p("</svg>\n"); }
 SVG &color(string nc) { return c = nc, *this; }
 void line(auto x1, auto y1, auto x2, auto y2) {
p("<line x1='$' y1='$' x2='$' y2='$' stroke='$'/>\n",
```

x1, -y1, x2, -y2, c); }

```
void circle(auto x, auto y, auto r) {
p("<circle cx='$' cy='$' r='$' stroke='$' "</pre>
                                                             void access(int x) {
   "fill='none'/>\n", x, -y, r, c); }
                                                              for (int u = x, last = 0; u; u = cur.pa) {
void text(auto x, auto y, string s, int w = 12) {
                                                               splay(u);
 p("<text x='$' y='$' font-size='$px'>$</text>\n",
                                                               cur.vir = cur.vir + o[rc].sub - o[last].sub;
  x, -y, w, s); }
                                                               rc = last; up(last = u);
                                                              splay(x);
2
     Data Structure
2.1 Dark Magic [095f25]
                                                             int find_root(int u) {
#include <ext/pb_ds/assoc_container.hpp>
                                                              int la = 0;
#include <ext/pb_ds/priority_queue.hpp>
                                                              for (access(u); u; u = lc) down(la = u);
using namespace __gnu_pbds;
                                                              return la;
// heap tags: paring/binary/binomial/rc_binomial/thin
template<typename T>
                                                             void split(int x, int y) { chroot(x); access(y); }
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>, \
                                                             void chroot(int u) { access(u); set_rev(u); }
                  pairing_heap_tag>;
                                                            public:
// pbds_heap::point_iterator
                                                             LCT(int n = 0) : o(n + 1) {}
  x = pq.push(10); pq.modify(x, 87); a.join(b);
                                                             int add(const Val &v = {}) {
// tree tags: rb_tree_tag/ov_tree_tag/splay_tree_tag
                                                              o.push_back(v);
template<typename T>
                                                              return int(o.size()) - 2;
using ordered_set = tree<T, null_type, less<T>,
   rb_tree_tag, tree_order_statistics_node_update>;
                                                             void set_val(int u, const Val &v) {
// find_by_order, order_of_key
                                                              splay(++u); cur.v = v; up(u);
// hash tables: cc_hash_table/gp_hash_table
2.2 Link-Cut Tree [60627f]
                                                             void set_sval(int u, const SVal &v) {
template <typename Val, typename SVal> class LCT {
                                                             access(++u); cur.sv = v; up(u);
struct node {
                                                             Val query(int x, int y) {
 int pa, ch[2];
 bool rev;
                                                              split(++x, ++y); return o[y].prod;
 Val v, prod, rprod;
 SVal sv, sub, vir;
                                                             SVal subtree(int p, int u) {
 node() : pa{0}, ch{0, 0}, rev{false}, v{}, prod{},
                                                              chroot(++p); access(++u);
                                                              return cur.vir + cur.sv;
    rprod{}, sv{}, sub{}, vir{} {};
#define cur o[u]
                                                             bool connected(int u, int v) {
#define lc cur.ch[0]
                                                              return find_root(++u) == find_root(++v); }
#define rc cur.ch[1]
                                                             void link(int x, int y) {
                                                              chroot(++x); access(++y);
vector<node> o:
                                                              o[y].vir = o[y].vir + o[x].sub;
bool is_root(int u) const {
 return o[cur.pa].ch[0] != u && o[cur.pa].ch[1] != u;
                                                              up(o[x].pa = y);
                                                             void cut(int x, int y) {
bool is_rch(int u) const {
                                                              split(++x, ++y);
 return o[cur.pa].ch[1] == u && !is_root(u); }
void down(int u) {
                                                              o[y].ch[0] = o[x].pa = 0; up(y);
 if (not cur.rev) return;
                                                            #undef cur
 if (lc) set_rev(lc);
                                                            #undef lc
 if (rc) set_rev(rc);
                                                            #undef rc
 cur.rev = false;
                                                           };
}
void up(int u) {
                                                            2.3 LiChao Segment Tree [b9c827]
 cur.prod = o[lc].prod * cur.v * o[rc].prod;
 cur.rprod = o[rc].rprod * cur.v * o[lc].rprod;
                                                            struct L {
 cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
                                                             int m, k, id;
                                                             L(): id(-1) {}
void set_rev(int u) {
                                                             L(int a, int b, int c) : m(a), k(b), id(c) {}
 swap(lc, rc), swap(cur.prod, cur.rprod);
                                                             int at(int x) { return m * x + k; }
 cur.rev ^= 1;
                                                            class LiChao {
void rotate(int u) {
                                                           private:
 int f = cur.pa, g = o[f].pa, l = is_rch(u);
if (cur.ch[l ^ 1]) o[cur.ch[l ^ 1]].pa = f;
                                                             int n; vector<L> nodes;
                                                             static int lc(int x) { return 2 * x + 1; }
 if (not is_root(f)) o[g].ch[is_rch(f)] = u;
                                                             static int rc(int x) { return 2 * x + 2; }
 o[f].ch[l] = cur.ch[l ^ 1];
                                                             void insert(int l, int r, int id, L ln) {
 cur.ch[l ^ 1] = f;
                                                              int m = (l + r) >> 1;
 cur.pa = g, o[f].pa = u;
                                                              if (nodes[id].id == -1)
 up(f);
                                                               return nodes[id] = ln, void();
                                                              bool atLeft = nodes[id].at(l) < ln.at(l);</pre>
                                                              if (nodes[id].at(m) < ln.at(m))</pre>
void splay(int u) {
 vector<int> stk = {u};
                                                               atLeft ^= 1, swap(nodes[id], ln);
                                                              if (r - l == 1) return;
 while (not is_root(stk.back()))
                                                              if (atLeft) insert(l, m, lc(id), ln);
   stk.push_back(o[stk.back()].pa);
 while (not stk.empty()) {
                                                              else insert(m, r, rc(id), ln);
  down(stk.back());
  stk.pop_back();
                                                             int query(int l, int r, int id, int x) {
                                                              int m = (l + r) >> 1, ret = 0;
                                                              if (nodes[id].id != -1) ret = nodes[id].at(x);
 for (int f = cur.pa; not is_root(u); f = cur.pa) {
  if (!is_root(f))
                                                              if (r - l == 1) return ret;
   rotate(is_rch(u) == is_rch(f) ? f : u);
                                                              if (x < m) return max(ret, query(l, m, lc(id), x));</pre>
  rotate(u);
                                                              return max(ret, query(m, r, rc(id), x));
 up(u);
                                                           public:
```

```
LiChao(int n_{-}): n(n_{-}), nodes(n * 4) {}
 void insert(L ln) { insert(0, n, 0, ln); }
                                                                 return l + 1 - sz;
 int query(int x) { return query(0, n, 0, x); }
                                                                } else sum = s, l++;
                                                               } while (lowbit(l) != l);
};
                                                               return n + 1:
2.4
      Treap [ae576c]
__gnu_cxx::sfmt19937 rnd(7122); // <ext/random>
                                                              int find_last(int r, auto &&check) {
namespace Treap {
                                                               if (r <= 0) return -1;
struct node {
                                                               r += sz; push(r - 1); Monoid sum; // identity
int size, pri; node *lc, *rc, *pa;
node() : size(1), pri(rnd()), lc(0), rc(0), pa(0) {}
                                                               do {
 void pull() {
                                                                while (r > 1 and (r & 1)) r >>= 1;
 size = 1; pa = 0;
                                                                if (auto s = nd[r] + sum; check(s)) {
  if (lc) { size += lc->size; lc->pa = this; }
                                                                 while (r < sz) {</pre>
 if (rc) { size += rc->size; rc->pa = this; }
                                                                  prop(r); r = (r << 1) | 1;
 }
                                                                   if (auto nxt = nd[r] + sum; not check(nxt))
                                                                    sum = nxt, r--;
int SZ(node *x) { return x ? x->size : 0; }
node *merge(node *L, node *R) {
  if (not L or not R) return L ? L : R;
                                                                 return r - sz;
                                                                } else sum = s;
 if (L->pri > R->pri)
                                                               } while (lowbit(r) != r);
  return L->rc = merge(L->rc, R), L->pull(), L;
                                                               return -1;
  return R->lc = merge(L, R->lc), R->pull(), R;
                                                                   Graph
void splitBySize(node *o, int k, node *&L, node *&R) {
                                                              3.1
                                                                   2-SAT (SCC) [09167a]
 if (not o) L = R = 0;
                                                              class TwoSat { // test @ CSES Giant Pizza
 else if (int s = SZ(o->lc) + 1; s <= k)
  L=o, splitBySize(o->rc, k-s, L->rc, R), L->pull();
                                                               int n; vector<vector<int>> G, rG, sccs;
 else
                                                               vector<int> ord, idx, vis, res;
  R=o, splitBySize(o->lc, k, L, R->lc), R->pull();
                                                               void dfs(int u) {
\frac{1}{N} = \frac{1}{N} \sum_{i=1}^{N} \frac{1}{N} \sum_{i=1}^{N} \frac{1}{N} 
                                                                vis[u] = true;
int getRank(node *o) { // 1-base
                                                                for (int v : G[u]) if (!vis[v]) dfs(v);
 int r = SZ(o->lc) + 1;
                                                                ord.push_back(u);
 for (; o->pa; o = o->pa)
  if (o->pa->rc == o) r += SZ(o->pa->lc) + 1;
                                                               void rdfs(int u) {
 return r;
                                                                vis[u] = false; idx[u] = sccs.size() - 1;
                                                                sccs.back().push_back(u);
} // namespace Treap
                                                                for (int v : rG[u]) if (vis[v]) rdfs(v);
2.5 Linear Basis [138d5d]
                                                               }
template <int BITS, typename S = int> struct Basis {
                                                              public:
 static constexpr S MIN = numeric_limits<S>::min();
                                                               TwoSat(int n_{-}): n(n_{-}), G(n), rG(n), idx(n), vis(n),
 array<pair<llu, S>, BITS> b;
                                                                   res(n) {}
 Basis() { b.fill({0, MIN}); }
                                                               void add_edge(int u, int v) {
 void add(llu x, S p) {
                                                                G[u].push_back(v); rG[v].push_back(u);
  for (int i = BITS-1; i>=0; i--) if (x >> i & 1) {
                                                               void orr(int x, int y) {
  if ((x ^ y) == 1) return;
   if (b[i].first == 0) return b[i]={x, p}, void();
   if (b[i].second < p)</pre>
                                                                add_edge(x ^ 1, y); add_edge(y ^ 1, x);
    swap(b[i].first, x), swap(b[i].second, p);
   x ^= b[i].first;
                                                               bool solve() {
  }
                                                                for (int i = 0; i < n; ++i) if (not vis[i]) dfs(i);</pre>
 optional<llu> query_kth(llu v, llu k) {
                                                                for (int u : ord | views::reverse)
                                                                 if (vis[u]) sccs.emplace_back(), rdfs(u);
  vector<pair<llu, int>> o;
                                                                for (int i = 0; i < n; i += 2)
  if (idx[i] == idx[i + 1]) return false;</pre>
  for (int i = 0; i < BITS; i++)</pre>
   if (b[i].first) o.emplace_back(b[i].first, i);
  if (k >= (1ULL << o.size())) return {};</pre>
                                                                vector<bool> c(sccs.size());
  for (int i = int(o.size()) - 1; i >= 0; i--)
                                                                for (size_t i = 0; i < sccs.size(); ++i)</pre>
   if ((k >> i & 1) ^ (v >> o[i].second & 1))
                                                                 for (int z : sccs[i])
    v ^= o[i].first;
                                                                   res[z] = c[i], c[idx[z ^ 1]] = !c[i];
  return v:
                                                                return true:
 Basis filter(S l) {
                                                               bool get(int x) { return res[x]; }
                                                               int get_id(int x) { return idx[x]; }
 Basis res = *this:
  for (int i = 0; i < BITS; i++)</pre>
                                                               int count() { return sccs.size(); }
   if (res.b[i].second < l) res.b[i] = {0, MIN};</pre>
  return res;
                                                              3.2
                                                                     BCC [6ac6db]
                                                              class BCC {
};
                                                               int n, ecnt, bcnt;
2.6
     Binary Search On Segtree [6c61c0]
                                                               vector<vector<pair<int, int>>> g;
// find_first = l \rightarrow minimal \times s.t. check([l, x))
                                                               vector<int> dfn, low, bcc, stk;
// find_last = r -> maximal x s.t. check( [x, r) )
                                                               vector<bool> ap, bridge;
int find_first(int l, auto &&check) {
                                                               void dfs(int u, int f)
 if (l >= n) return n + 1;
                                                                dfn[u] = low[u] = dfn[f] + 1;
 l += sz; push(l); Monoid sum; // identity
                                                                int ch = 0;
                                                                for (auto [v, t] : g[u]) if (bcc[t] == -1) {
  while ((l & 1) == 0) l >>= 1;
                                                                 bcc[t] = 0; stk.push_back(t);
  if (auto s = sum + nd[l]; check(s)) {
                                                                 if (dfn[v]) {
   while (l < sz) {</pre>
                                                                  low[u] = min(low[u], dfn[v]);
    prop(l); l = (l << 1);
                                                                  continue;
    if (auto nxt = sum + nd[l]; not check(nxt))
                                                                 ++ch, dfs(v, u);
     sum = nxt, l++;
```

```
low[u] = min(low[u], low[v]);
if (low[v] > dfn[u]) bridge[t] = true;
                                                               for (int i = 0; i < n; i++)</pre>
   if (low[v] < dfn[u]) continue;</pre>
                                                               if (in[i] == -1) dfs(dfs, i, -1);
                                                               for (int i = 0; i < n; i++)
  if (dsu.anc(i) == i) id[i] = cnt++;</pre>
   ap[u] = true;
   while (not stk.empty()) {
    int o = stk.back(); stk.pop_back();
                                                               vector<vector<int>> comps(cnt);
    bcc[o] = bcnt;
                                                               for (int i = 0; i < n; i++)</pre>
                                                               comps[id[dsu.anc(i)]].push_back(i);
    if (o == t) break;
                                                               return comps;
   bcnt += 1;
                                                             } // test @ yosupo judge
  }
                                                              3.5 DMST [f4317e]
  ap[u] = ap[u] and (ch != 1 or u != f);
                                                             using lld = int64_t;
                                                             struct E { int s, t; lld w; }; // 0-base
public:
                                                             struct PQ {
 BCC(int n_{-}) : n(n_{-}), ecnt(0), bcnt(0), g(n), dfn(n),
                                                               struct P {
    low(n), stk(), ap(n) {}
                                                               lld v; int i;
 void add_edge(int u, int v) {
                                                                bool operator>(const P &b) const { return v > b.v; }
  g[u].emplace_back(v, ecnt);
  g[v].emplace_back(u, ecnt++);
                                                               min_heap<P> pq; lld tag;
                                                               void push(P p) { p.v -= tag; pq.emplace(p); }
 void solve() {
                                                               P top() { P p = pq.top(); p.v += tag; return p; }
  bridge.assign(ecnt, false); bcc.assign(ecnt, -1);
                                                               void join(PQ &b) {
  for (int i = 0; i < n; ++i) if (!dfn[i]) dfs(i, i);</pre>
                                                                if (pq.size() < b.pq.size())</pre>
                                                                 swap(pq, b.pq), swap(tag, b.tag);
 int bcc_id(int x) const { return bcc[x]; }
                                                                while (!b.pq.empty()) push(b.top()), b.pq.pop();
 bool is_ap(int x) const { return ap[x];
bool is_bridge(int x) const { return bridge[x]; }
                                                             };
                                                             vector<int> dmst(const vector<E> &e, int n, int root) {
3.3
      Round Square Tree [cf6d74]
                                                               vector<PQ> h(n * 2);
struct RST { // be careful about isolate point
                                                               for (int i = 0; i < int(e.size()); ++i)</pre>
 int n; vector<vector<int>> T;
                                                               h[e[i].t].push({e[i].w, i});
                                                               vector<int> a(n * 2); iota(all(a), 0);
vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
 RST(auto &G) : n(int(G.size())), T(n) {
  vector<int> stk, vis(n), low(n);
                                                               auto o = [\&](auto Y, int X) \rightarrow int {
  auto dfs = [&](auto self, int u, int d) -> void {
   low[u] = vis[u] = d; stk.push_back(u);
                                                               return x==a[x] ? x : a[x] = Y(Y, a[x]); };
   for (int v : G[u]) if (!vis[v]) {
                                                               auto S = [&](int i) { return o(o, e[i].s); };
    self(self, v, d + 1);
                                                               int pc = v[root] = n;
    if (low[v] = vis[u]) {
                                                               for (int i = 0; i < n; ++i) if (v[i] == -1)
     int cnt = int(T.size()); T.emplace_back();
                                                                for (int p = i; v[p]<0 || v[p]==i; p = S(r[p])) {</pre>
     for (int x = -1; x != v; stk.pop_back())
                                                                 if (v[p] == i)
      T[cnt].push_back(x = stk.back());
                                                                  for (int q = pc++; p != q; p = S(r[p])) {
     T[u].push_back(cnt); // T is rooted
                                                                   h[p].tag -= h[p].top().v; h[q].join(h[p]);
    } else low[u] = min(low[u], low[v]);
                                                                   pa[p] = a[p] = q;
   } else low[u] = min(low[u], vis[v]);
  };
                                                                 while (S(h[p].top().i) == p) h[p].pq.pop();
  for (int u = 0; u < n; u++)</pre>
                                                                v[p] = i; r[p] = h[p].top().i;
   if (!vis[u]) dfs(dfs, u, 1);
 } // T may be forest; after dfs, stk are the roots
                                                               vector<int> ans;
                                                               for (int i = pc - 1; i >= 0; i--) if (v[i] != n) {
}; // test @ 2020 Shanghai K
                                                                for (int f = e[r[i]].t; f!=-1 && v[f]!=n; f = pa[f])
3.4 Edge TCC [5a2668]
                                                                v[f] = n;
vector<vector<int>> ETCC(auto &adj) {
                                                               ans.push_back(r[i]);
 const int n = static_cast<int>(adj.size());
 vector<int> up(n), low(n), in, out, nx, id;
                                                              return ans; // default minimize, returns edgeid array
 in = out = nx = id = vector<int>(n, -1);
 int dfc = 0, cnt = 0; Dsu dsu(n);
                                                              3.6 Dominator Tree [ea5b7c]
 auto merge = [&](int u, int v) {
  dsu.join(u, v); up[u] += up[v]; };
                                                             struct Dominator {
                                                              vector<vector<int>> g, r, rdom; int tk;
vector<int> dfn, rev, fa, sdom, dom, val, rp;
 auto dfs = [&](auto self, int u, int p) -> void {
  in[u] = low[u] = dfc++;
  for (int v : adj[u]) if (v != u) {
                                                               Dominator(int n) : g(n), r(n), rdom(n), tk(0) {
   if (v == p) { p = -1; continue; }
                                                               dfn = rev = fa = sdom = dom =
   if (in[v] == -1) {
                                                                 val = rp = vector<int>(n, -1); }
                                                               void add_edge(int x, int y) { g[x].push_back(y); }
    self(self, v, u);
    if (nx[v] == -1 && up[v] <= 1) {
                                                               void dfs(int x) {
     up[u] += up[v]; low[u] = min(low[u], low[v]);
                                                               rev[dfn[x] = tk] = x;
     continue;
                                                               fa[tk] = sdom[tk] = val[tk] = tk; tk++;
                                                               for (int u : g[x]) {
   if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
    if (up[v] == 0) v = nx[v];
    if (low[u] > low[v])
                                                                 r[dfn[u]].push_back(dfn[x]);
     low[u] = low[v], swap(nx[u], v);
                                                               }
    for (; v != -1; v = nx[v]) merge(u, v);
   } else if (in[v] < in[u]) {</pre>
                                                               void merge(int x, int y) { fa[x] = y; }
                                                               int find(int x, int c = 0) {
    low[u] = min(low[u], in[v]); up[u]++;
                                                                if (fa[x] == x) return c ? -1 : x;
    for (int &x = nx[u]; x != -1 &&
                                                                if (int p = find(fa[x], 1); p != -1) {
      in[x] \le in[v] \&\& in[v] \le out[x]; x = nx[x])
                                                                 if (sdom[val[x]] > sdom[val[fa[x]]])
     merge(u, x);
                                                                  val[x] = val[fa[x]];
    up[u]--;
                                                                 fa[x] = p;
                                                                 return c ? p : val[x];
   }
                                                                } else return c ? fa[x] : val[x];
  out[u] = dfc;
```

vector<int64_t> Sub, Sub2;

vector<int> Cnt, Cnt2;

```
vector<int> build(int s, int n) {
  // return the father of each node in dominator tree
                                                            void DfsSz(vector<int> &tmp, int x) {
                                                             vis[x] = true, sz[x] = 1, mx[x] = 0;
  dfs(s); // p[i] = -2 if i is unreachable from s
                                                             for (auto [u, w] : g[x]) if (not vis[u]) {
  for (int i = tk - 1; i >= 0; --i) {
                                                              DfsSz(tmp, u); sz[x] += sz[u];
   for (int u : r[i])
                                                              mx[x] = max(mx[x], sz[u]);
    sdom[i] = min(sdom[i], sdom[find(u)]);
   if (i) rdom[sdom[i]].push_back(i);
                                                             tmp.push_back(x);
   for (int u : rdom[i]) {
    int p = find(u);
                                                            void DfsDist(int x, int64_t D = 0) {
   dom[u] = (sdom[p] == i ? i : p);
                                                             Dist[x].push_back(D); vis[x] = true;
                                                             for (auto [u, w] : g[x])
                                                              if (not vis[u]) DfsDist(u, D + w);
   if (i) merge(i, rp[i]);
  }
  vector<int> p(n, -2); p[s] = -1;
                                                            void DfsCen(int x, int D, int p) {
                                                             vector<int> tmp; DfsSz(tmp, x);
 for (int i = 1; i < tk; ++i)</pre>
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
                                                             int M = int(tmp.size()), C = -1;
  for (int i = 1; i < tk; ++i)</pre>
                                                             for (int u : tmp)
                                                              if (max(M - sz[u], mx[u]) * 2 <= M) C = u;</pre>
  p[rev[i]] = rev[dom[i]];
                                                             for (int u : tmp) vis[u] = false;
 return p;
} // test @ yosupo judge
                                                             DfsDist(C);
                                                             for (int u : tmp) vis[u] = false;
};
                                                             pa[C] = p, vis[C] = true, dep[C] = D;
3.7
      Edge Coloring [029763]
                                                             for (auto [u, w] : g[C])
// max(d_u) + 1 edge coloring, time: O(NM)
                                                              if (not vis[u]) DfsCen(u, D + 1, C);
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
                                                           public:
for (int i = 0; i <= N; i++)</pre>
                                                            Centroid(int N) : g(N), pa(N), dep(N),
 for (int j = 0; j <= N; j++)
                                                             vis(N), sz(N), mx(N), Dist(N),
    C[i][j] = G[i][j] = 0;
                                                             Sub(N), Sub2(N), Cnt(N), Cnt2(N) {}
                                                            void AddEdge(int u, int v, int w) {
void solve(vector<pair<int, int>> &E, int N) {
                                                             g[u].emplace_back(v, w);
int X[kN] = {}, a;
                                                             g[v].emplace_back(u, w);
auto update = [&](int u) {
 for (X[u] = 1; C[u][X[u]]; X[u]++);
                                                            void Build() { DfsCen(0, 0, -1); }
                                                            void Mark(int v) {
auto color = [&](int u, int v, int c) {
                                                             int x = v, z = -1;
 int p = G[u][v];
                                                             for (int i = dep[v]; i >= 0; --i) {
 G[u][v] = G[v][u] = c;
                                                              Sub[x] += Dist[v][i], Cnt[x]++;
 C[u][c] = v, C[v][c] = u;
                                                              if (z != -1)
 C[u][p] = C[v][p] = 0;
                                                               Sub2[z] += Dist[v][i], Cnt2[z]++;
 if (p) X[u] = X[v] = p;
                                                              x = pa[z = x];
 else update(u), update(v);
  return p;
};
                                                            int64_t Query(int v) {
auto flip = [&](int u, int c1, int c2) {
                                                             int64_t res = 0;
 int p = C[u][c1];
                                                             int x = v, z = -1
 swap(C[u][c1], C[u][c2]);
                                                             for (int i = dep[v]; i >= 0; --i) {
 if (p) G[u][p] = G[p][u] = c2;
                                                              res += Sub[x] + 1LL * Cnt[x] * Dist[v][i];
  if (!C[u][c1]) X[u] = c1;
                                                              if (z != -1)
 if (!C[u][c2]) X[u] = c2;
                                                               res -= Sub2[z] + 1LL * Cnt2[z] * Dist[v][i];
  return p;
                                                              x = pa[z = x];
};
for (int i = 1; i <= N; i++) X[i] = 1;
                                                             return res;
 for (int t = 0; t < E.size(); t++) {</pre>
 auto [u, v] = E[t];
                                                           }; // pa, dep are centroid tree attributes
  int v0 = v, c = X[u], c0 = c, d;
 vector<pair<int, int>> L; int vst[kN] = {};
                                                           3.9 Lowbit Decomposition [did724]
 while (!G[u][v0]) {
                                                           class LBD {
   L.emplace_back(v, d = X[v]);
                                                            int timer, chains;
   if (!C[v][c]) for (a=L.size()-1;a>=0;a--)
                                                            vector<vector<int>> G;
     c = color(u, L[a].first, c);
                                                            vector<int> tl, tr, chain, head, dep, pa;
   else if (!C[u][d]) for (a=L.size()-1;a>=0;a--)
                                                            // chains : number of chain
    color(u, L[a].first, L[a].second);
                                                            // tl, tr[u] : subtree interval in the seq. of u
   else if (vst[d]) break;
                                                            // head[i] : head of the chain i
   else vst[d] = 1, v = C[u][d];
                                                            // chian[u] : chain id of the chain u is on
                                                            void predfs(int u, int f) {
 if (!G[u][v0]) {
                                                             dep[u] = dep[pa[u] = f] + 1;
   for (; v; v = flip(v, c, d), swap(c, d));
                                                             for (int v : G[u]) if (v != f) {
   if (C[u][c0]) { a = int(L.size()) - 1;
                                                              predfs(v, u):
   while (--a >= 0 && L[a].second != c);
                                                              if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
                                                               chain[u] = chain[v];
   } else t--;
                                                             if (chain[u] == 0) chain[u] = ++chains;
}
                                                            void dfschain(int u, int f) {
      Centroid Decomposition [670cdd]
                                                             tl[u] = timer++;
                                                             if (head[chain[u]] == -1)
class Centroid {
                                                              head[chain[u]] = u;
vector<vector<pair<int, int>>> g; // g[u] = {(v, w)}
vector<int> pa, dep, vis, sz, mx;
                                                             for (int v : G[u])
                                                              if (v != f and chain[v] == chain[u])
vector<vector<int64_t>> Dist;
```

dfschain(v, u);

for (int v : G[u])

```
3.14
                                                                           MaximalClique [293730]
   if (v != f and chain[v] != chain[u])
    dfschain(v, u);
                                                                  // contain a self loop u to u, than u won't in clique
                                                                  template <size_t maxn> class MaxClique {
  tr[u] = timer;
                                                                  private:
                                                                    using bits = bitset<maxn>;
public:
 LBD(auto &&G_) : n((int)size(G_)),
                                                                    bits popped, G[maxn], ans;
  \begin{array}{l} \mbox{timer}(0)\,,\; \mbox{chains}(0)\,,\; \mbox{G}(\mbox{G}_{\_}),\; \mbox{tl}(\mbox{n})\,,\; \mbox{tr}(\mbox{n})\,,\\ \mbox{chain}(\mbox{n})\,,\; \mbox{head}(\mbox{n}\,+\,1,\,-1)\,,\; \mbox{dep}(\mbox{n})\,,\; \mbox{pa}(\mbox{n}) \end{array}
                                                                    size_t deg[maxn], deo[maxn], n;
                                                                    void sort_by_degree() {
                                                                     popped.reset();
  { predfs(0, 0); dfschain(0, 0); }
                                                                     for (size_t i = 0; i < n; ++i)</pre>
 PII get_subtree(int u) { return {tl[u], tr[u]}; }
                                                                      deg[i] = G[i].count();
 vector<PII> get_path(int u, int v) {
  vector<PII> res;
                                                                     for (size_t i = 0; i < n; ++i) {</pre>
                                                                      size_t mi = maxn, id = 0;
  while (chain[u] != chain[v]) {
                                                                      for (size_t j = 0; j < n; ++j)</pre>
   if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
                                                                       if (not popped[j] and deg[j] < mi)</pre>
    swap(u, v);
                                                                        mi = deg[id = j];
   int s = head[chain[u]];
                                                                      popped[deo[i] = id] = 1;
   res.emplace_back(tl[s], tl[u] + 1);
                                                                      for (size_t u = G[i]._Find_first(); u < n;</pre>
   u = pa[s];
                                                                        u = G[i]._Find_next(u))
  if (dep[u] < dep[v]) swap(u, v);</pre>
                                                                       --deg[u];
 res.emplace_back(tl[v], tl[u] + 1);
  return res;
                                                                    void BK(bits R, bits P, bits X) {
                                                                     if (R.count() + P.count() <= ans.count()) return;</pre>
}; // 記得在資結上對點的修改要改成對其 dfs 序的修改
                                                                     if (not P.count() and not X.count()) {
3.10 Virtual Tree [ad5cf5]
                                                                      if (R.count() > ans.count()) ans = R;
vector<pair<int, int>> build(vector<int> vs, int r) {
                                                                      return;
 vector<pair<int, int>> res;
 sort(vs.begin(), vs.end(), [](int i, int j) {
  return dfn[i] < dfn[j]; });</pre>
                                                                     /* greedily chosse max degree as pivot
bits cur = P | X; size_t pivot = 0, sz = 0;
 vector<int> s = {r};
                                                                     for ( size_t u = cur._Find_first() ;
 for (int v : vs) if (v != r) {
                                                                     u < n ; u = cur._Find_next( u ) )
if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
  if (int o = lca(v, s.back()); o != s.back()) {
   while (s.size() >= 2) {
                                                                     cur = P & ( ~G[ pivot ] );
    if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
                                                                     */ // or simply choose first
    res.emplace_back(s[s.size() - 2], s.back());
                                                                     bits cur = P & (~G[(P | X)._Find_first()]);
    s.pop_back();
                                                                     for (size_t u = cur._Find_first(); u < n;</pre>
                                                                       u = cur._Find_next(u)) {
   if (s.back() != o) {
                                                                      if (R[u]) continue;
    res.emplace_back(o, s.back());
                                                                      R[u] = 1;
    s.back() = o;
                                                                      BK(R, P & G[u], X & G[u]);
                                                                      R[u] = P[u] = 0, X[u] = 1;
  s.push_back(v);
                                                                   }
                                                                  public:
 for (size_t i = 1; i < s.size(); ++i)</pre>
                                                                   void init(size_t n_) {
  res.emplace_back(s[i - 1], s[i]);
                                                                     n = n;
 return res; // (x, y): x->y
                                                                     for (size_t i = 0; i < n; ++i) G[i].reset();</pre>
                                                                     ans.reset();
3.11 Tree Hashing [d6a9f9]
vector<int> g[maxn]; llu h[maxn];
                                                                   void add_edges(int u, bits S) { G[u] = S; }
void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
llu F(llu z) { // xorshift64star from iwiwi
 z \stackrel{\wedge}{=} z >> 12; z \stackrel{\wedge}{=} z << 25; z \stackrel{\wedge}{=} z >> 27;
                                                                    int solve() {
return z * 2685821657736338717LL;
                                                                     sort_by_degree(); // or simply iota( deo... )
for (size_t i = 0; i < n; ++i)</pre>
llu hsah(int u, int f) {
                                                                      deg[i] = G[i].count();
llu r = 127; // bigger?
                                                                     bits pob, nob = 0; pob.set();
 for (int v : g[u]) if (v != f) r += hsah(v, u);
                                                                     for (size_t i = n; i < maxn; ++i) pob[i] = 0;
for (size_t i = 0; i < n; ++i) {</pre>
 return h[u] = F(r);
} // test @ UOJ 763 & yosupo library checker
                                                                      size_t v = deo[i];
3.12 Mo's Algorithm on Tree
                                                                      bits tmp;
dfs u:
                                                                      tmp[v] = 1;
push u
                                                                      BK(tmp, pob & G[v], nob & G[v]);
 iterate subtree
                                                                      pob[v] = 0, nob[v] = 1;
push u
Let P = LCA(u, v) with St(u) \le St(v)
                                                                     return static_cast<int>(ans.count());
if (P == u) query[St(u), St(v)]
                                                                   }
else query[Ed(u), St(v)], query[St(P), St(P)]
                                                                         MaximumClique [aee5d8]
3.13
       Count Cycles [c7e8f2]
// ord = sort by deg decreasing, rk[ord[i]] = i
                                                                   constexpr size_t kN = 150; using bits = bitset<kN>;
// D[i] = edge point from rk small to rk big
                                                                   struct MaxClique {
for (int x : ord) { // c3
                                                                   bits G[kN], cs[kN];
 for (int y : D[x]) vis[y] = 1;
                                                                    int ans, sol[kN], q, cur[kN], d[kN], n;
                                                                    void init(int _n) {
 for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
                                                                     n = _n;
 for (int y : D[x]) vis[y] = 0;
                                                                     for (int i = 0; i < n; ++i) G[i].reset();</pre>
for (int x : ord) { // c4
  for (int y : D[x]) for (int z : adj[y])
                                                                    void add_edge(int u, int v) { G[u][v] = G[v][u] = 1; }
  if (rk[z] > rk[x]) c4 += vis[z]++;
                                                                    void pre_dfs(vector<int> &v, int i, bits mask) {
 for (int y : D[x]) for (int z : adj[y])
                                                                     if (i < 4) {
  if (rk[z] > rk[x]) --vis[z];
                                                                      for (int x : v) d[x] = (int)(G[x] \& mask).count();
                                                                      sort(all(v), [&](int x, int y) {
} // both are O(M*sqrt(M)), test @ 2022 CCPC guangzhou
```

```
return d[x] > d[y]; });
  vector<int> c(v.size());
  cs[1].reset(), cs[2].reset();
  int l = max(ans - q + 1, 1), r = 2, tp = 0, k;
  for (int p : v) {
   for (k = 1; (cs[k] & G[p]).any(); ++k);
   if (k >= r) cs[++r].reset();
   cs[k][p] = 1;
   if (k < l) v[tp++] = p;
 for (k = l; k < r; ++k)
for (auto p = cs[k]._Find_first();</pre>
     p < kN; p = cs[k]._Find_next(p))
    v[tp] = (int)p, c[tp] = k, ++tp;
 dfs(v, c, i + 1, mask);
void dfs(vector<int> &v, vector<int> &c,
   int i, bits mask) {
  while (!v.empty()) {
   int p = v.back(); v.pop_back(); mask[p] = 0;
   if (q + c.back() <= ans) return;</pre>
                                                                 };
   cur[q++] = p;
   vector<int> nr;
   for (int x : v) if (G[p][x]) nr.push_back(x);
   if (!nr.empty()) pre_dfs(nr, i, mask & G[p]);
   else if (q > ans) ans = q, copy_n(cur, q, sol);
   c.pop_back(); --q;
 }
int solve() {
 vector<int> v(n); iota(all(v), 0);
  ans = q = 0; pre_dfs(v, 0, bits(string(n, '1')));
  return ans; // sol[0 ~ ans-1]
} cliq; // test @ yosupo judge
3.16
      Minimum Mean Cycle [e23bc0]
// WARNING: TYPE matters
struct Edge { int s, t; llf c; };
llf solve(vector<Edge> &e, int n) {
  // O(VE), returns inf if no cycle, mmc otherwise
vector<VI> prv(n + 1, VI(n)), prve = prv;
vector<vector<llf>>> d(n + 1, vector<llf>(n, inf));
d[0] = vector<llf>(n, 0);
for (int i = 0; i < n; i++) {</pre>
  for (int j = 0; j < (int)e.size(); j++) {</pre>
   auto [s, t, c] = e[j];
   if (d[i][s] < inf && d[i + 1][t] > d[i][s] + c) {
    d[i + 1][t] = d[i][s] + c;
    prv[i + 1][t] = s; prve[i + 1][t] = j;
 }
llf mmc = inf; int st = -1;
                                                                 };
for (int i = 0; i < n; i++) {</pre>
 llf avg = -inf;
 for (int k = 0; k < n; k++) {
  if (d[n][i] < inf - eps)</pre>
    avg = max(avg, (d[n][i] - d[k][i]) / (n - k));
   else avg = inf;
  if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
if (st == -1) return inf;
 vector<int> vst(n), eid, cycle, rho;
 for (int i = n; !vst[st]; st = prv[i--][st]) {
 vst[st]++; eid.emplace_back(prve[i][st]);
  rho.emplace_back(st);
while (vst[st] != 2) {
 int v = rho.back(); rho.pop_back();
                                                                   Construct minimum vertex cover from maximum matching {\it M} on bipartite
  cycle.emplace_back(v); vst[v]++;
                                                                   1. Redirect every edge: y \to x if (x,y) \in M, x \to y otherwise. 2. DFS from unmatched vertices in X.
reverse(all(eid)); eid.resize(cycle.size());
return mmc;
                                                                 · Minimum cost cyclic flow
4
     Flow & Matching
                                                                   2. For each edge (x,y,c), connect x \to y with (cost,cap)=(c,1) if c>0, otherwise connect y \to x with (cost,cap)=(-c,1)
4.1 HopcroftKarp [930040]
struct HK {
vector<int> l, r, a, p; int ans;
                                                                   4. For each vertex v with d(v)>0, connect S\to v with (cost,cap)=
```

 $HK(int n, int m, auto \&g) : l(n,-1), r(m,-1), ans(0) {$

if (l[i] == -1) q.push(a[i] = p[i] = i); // bitset<maxn> nvis, t; nvis.set(); while (!q.empty()) { int z, x = q.front(); q.pop(); if (l[a[x]] != -1) continue; for (int y : g[x]) { // or iterate t = g[x]&nvis // nvis.reset(y); **if** (r[y] == -1) { for (z = y; z != -1;)r[z] = x, swap(l[x], z), x = p[x]; match = true; ++ans; break; } else if (p[r[y]] == -1)q.push(z = r[y]), p[z] = x, a[z] = a[x]; } } } Kuhn Munkres [2c09ed] 4.2 struct KM { // maximize, test @ UOJ 80 int n, l, r; lld ans; // fl and fr are the match vector<lld> hl, hr; vector<int> fl, fr, pre, q; void bfs(const auto &w, int s) { vector<int> vl(n), vr(n); vector<lld> slk(n, INF); l = r = 0; vr[q[r++] = s] = **true**; const auto check = [&](int x) -> bool { if (vl[x] || slk[x] > 0) return true; vl[x] = true; slk[x] = INF; if (fl[x] != -1) return vr[q[r++] = fl[x]] = true; while (x != -1) swap(x, fr[fl[x] = pre[x]]);return false: }; while (true) { while (l < r)for (int x = 0, y = q[l++]; x < n; ++x) if (!vl[x])</pre> if (chmin(slk[x], hl[x] + hr[y] - w[x][y]))if (pre[x] = y, !check(x)) return; lld d = ranges::min(slk); for (int x = 0; x < n; ++x) vl[x] ? hl[x] += d : slk[x] -= d; for (int x = 0; x < n; ++x) if (vr[x]) hr[x] -= d; for (int x = 0; x < n; ++x) if (!check(x)) return; $KM(int n_{, const auto \&w) : n(n_{, ans(0), ans(0),$ hl(n), hr(n), fl(n, -1), fr(fl), pre(n), q(n) { for (int i = 0; i < n; ++i) hl[i]=ranges::max(w[i]); for (int i = 0; i < n; ++i) bfs(w, i);</pre> for (int i = 0; i < n; ++i) ans += w[i][fl[i]];</pre> 4.3 Flow Models Maximum/Minimum flow with lower bound / Circulation problem 1. Construct super source S and sink T. For each edge (x, y, l, u), connect $x \to y$ with capacity u - lFor each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds. 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v). To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer. Also, f is a mincost valid flow. To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer. 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.

3. $x \in X$ is chosen iff x is unvisited; $y \in Y$ is chosen iff y is visited.

For each edge with c<0, sum these cost as K, then increase d(y) by 1,

1. Consruct super source S and sink T

decrease d(x) by 1

(0, d(v))

for (bool match = true; match;) {

for (int i = 0; i < n; i++)</pre>

queue<int> q;

match = false; a.assign(n, -1); p.assign(n, -1);

```
5. For each vertex v with d(v)~<~ 0, connect v~\rightarrow~T with (cost, cap)~=~
                                                                                                            struct Edge { int to, rev; T flow, cap; };
                                                                                                            int n, mx; vector<vector<Edge>> adj; vector<T> excess;
    6. Flow from S to T, the answer is the cost of the flow C+K
                                                                                                            vector<int> d, cnt, active; vector<vector<int>> B;
 · Maximum density induced subgraph
                                                                                                            void add_edge(int u, int v, int f) {
Edge a{v, (int)size(adj[v]), 0, f};
    1. Binary search on answer, suppose we're checking answer T
    2. Construct a max flow model, let K be the sum of all weights
                                                                                                             Edge b{u, (int)size(adj[u]), 0, 0};
    3. Connect source s \rightarrow v, v \in G with capacity K
        For each edge (u,v,w) in G, connect u \to v and v \to u with capacity w
                                                                                                             adj[u].push_back(a), adj[v].push_back(b);
    5. For v~\in~G, connect it with sink v~\to~t with capacity K~+~2T~-
         \left(\sum_{e \in E(v)} w(e)\right) - 2w(v)
                                                                                                            void enqueue(int v) {
                                                                                                             if (!active[v] && excess[v] > 0 && d[v] < n) {</pre>
    6. \hat{T} is a valid answer if the maximum flow f < K|V|
  Minimum weight edge cover
                                                                                                               mx = max(mx, d[v]);
    1. For each v \in V create a copy v', and connect u' \to v' with weight
                                                                                                               B[d[v]].push_back(v); active[v] = 1;
                                                                                                             }
    2. Connect v \to v' with weight 2\mu(v), where \mu(v) is the cost of the cheap-
        est edge incident to v.
   3. Find the minimum weight perfect matching on G'. Project selection cheat sheet: S,T 分別代表 0,1 側,最小化總花費。
                                                                                                            void push(int v, Edge &e) {
                                                                                                             T df = min(excess[v], e.cap - e.flow);
      i 為 O 時花費 c
                                               (i, T, c)
                                                                                                             if (df <= 0 || d[v] != d[e.to] + 1) return;</pre>
      i 為 1 時花費 c
                                               (S, i, c)
                                                                                                             e.flow += df, adj[e.to][e.rev].flow -= df;
      i\in I 有任何一個為 0 時花費 c i\in I 有任何一個為 1 時花費 c
                                               (i, w, \infty), (w, T, c)
                                                                                                             excess[e.to] += df, excess[v] -= df;
                                               (S, w, c), (w, i, \infty)
直接得到 c; (S, i, c)
      i 為 O 時得到 c
                                                                                                             enqueue(e.to);
      i 為 1 時得到 c i 為 0 ,j 為 1 時花費 c
                                               直接得到 c; (i, T, c)
                                               (i, j, c)
                                                                                                            void gap(int k) {
      i,j 不同時花費 c
                                               (i, j, c), (j, i, c)
                                                                                                             for (int v = 0; v < n; v++) if (d[v] >= k)
      i,j 同時是 0 時得到 c
                                               直接得到c; (S, w, c), (w, i, \infty), (w, j, \infty)
                                                                                                               cnt[d[v]]--, d[v] = n, cnt[d[v]]++;
      i,j 同時是 1 時得到 c
                                              直接得到 c ; (i, w, \infty), (j, w, \infty), (w, T, c)
   Submodular functions minimization  \begin{aligned} & \text{--Exp}(F, u, w, \infty), (j, w, \infty), (w, x, v_f) \\ & \text{--Exp}(F, u, x, w_f), (j, w, \infty), (w, x, v_f) \\ & \text{--Exp}(F, u, x, w_f) \\ & \text{--Exp}(F, u,
                                                                                                            void relabel(int v) {
                                                                                                             cnt[d[v]]--; d[v] = n;
                                                                                                             for (auto e : adj[v])
                                                                                                               if (e.cap > e.flow) d[v] = min(d[v], d[e.to] + 1);
                                                                                                             cnt[d[v]]++; enqueue(v);
   - Add edges (i, j, \phi_{ij}(0, 1) + \phi_{ij}(1, 0) - \phi_{ij}(0, 0) - \phi_{ij}(1, 1)).

- Denote x_{ijk} as helper nodes. Let P = \psi_{ijk}(0, 0, 0) + \psi_{ijk}(0, 1, 1) + \psi_{ijk}(1, 0, 1) + \psi_{ijk}(1, 1, 0) - \psi_{ijk}(0, 0, 1) - \psi_{ijk}(0, 1, 0) - \psi_{ijk}(1, 1, 1). Add -P to answer. If P \geq 0, add edges (i, x_{ijk}, P)
                                                                                                            void discharge(int v) {
                                                                                                             for (auto &e : adj[v])
                                                                                                               if (excess[v] > 0) push(v, e);
       (j,x_{ijk},P), (k,x_{ijk},P), (x_{ijk},T,P); otherwise (x_{ijk},i,-P), (x_{ijk},j,-P), (x_{ijk},k,-P). (S,x_{iik},-P).
                                                                                                               else break;
      (x_{ijk}, k, -P), (S, x_{ijk}, -P). The minimum cut of this graph will be the the minimum value of the
                                                                                                              if (excess[v] <= 0) return;</pre>
                                                                                                             if (cnt[d[v]] == 1) gap(d[v]);
       function above.
                                                                                                             else relabel(v);
            Dinic [32c53e]
 template <typename Cap = int64_t> class Dinic {
                                                                                                            T max_flow(int s, int t) {
 private:
                                                                                                             for (auto &e : adj[s]) excess[s] += e.cap;
   struct E { int to, rev; Cap cap; }; int n, st, ed;
                                                                                                             cnt[0] = n; enqueue(s); active[t] = 1;
   vector<vector<E>> G; vector<size_t> lv, idx;
                                                                                                             for (mx = 0; mx >= 0;)
   bool BFS(int k) {
                                                                                                               if (!B[mx].empty()) {
     lv.assign(n, 0); idx.assign(n, 0);
                                                                                                                 int v = B[mx].back(); B[mx].pop_back();
    queue<int> bfs; bfs.push(st); lv[st] = 1;
                                                                                                                 active[v] = 0; discharge(v);
    while (not bfs.empty() and not lv[ed]) {
                                                                                                               } else --mx:
       int u = bfs.front(); bfs.pop();
                                                                                                             return excess[t];
      for (auto e: G[u]) if (e.cap >> k and !lv[e.to])
        bfs.push(e.to), lv[e.to] = lv[u] + 1;
                                                                                                            HLPP(int _n) : n(_n), adj(n), excess(n),
                                                                                                             d(n), cnt(n + 1), active(n), B(n) {}
     return lv[ed];
                                                                                                          4.6 Global Min-Cut [ae7013]
   Cap DFS(int u, Cap f = numeric_limits<Cap>::max()) {
                                                                                                          void add_edge(auto &w, int u, int v, int c) {
    if (u == ed) return f;
                                                                                                            w[u][v] += c; w[v][u] += c; }
    Cap ret = 0;
                                                                                                          auto phase(const auto &w, int n, vector<int> id) {
     for (auto &i = idx[u]; i < G[u].size(); ++i) {</pre>
                                                                                                           vector<lld> g(n); int s = -1, t = -1;
       auto &[to, rev, cap] = G[u][i];
                                                                                                            while (!id.empty()) {
       if (cap <= 0 or lv[to] != lv[u] + 1) continue;</pre>
                                                                                                             int c = -1;
      Cap nf = DFS(to, min(f, cap));
ret += nf; cap -= nf; f -= nf;
                                                                                                             for (int i : id) if (c == -1 || g[i] > g[c]) c = i;
                                                                                                             s = t; t = c;
      G[to][rev].cap += nf;
                                                                                                             id.erase(ranges::find(id, c));
      if (f == 0) return ret;
                                                                                                             for (int i : id) g[i] += w[c][i];
     if (ret == 0) lv[u] = 0;
                                                                                                            return tuple{s, t, g[t]};
     return ret;
                                                                                                          lld mincut(auto w, int n) {
   lld cut = numeric_limits<lld>::max();
 public:
   void init(int n_) { G.assign(n = n_, vector<E>()); }
                                                                                                            vector<int> id(n); iota(all(id), 0);
   void add_edge(int u, int v, Cap c) {
                                                                                                            for (int i = 0; i < n - 1; ++i) {
    G[u].push_back({v, int(G[v].size()), c});
                                                                                                             auto [s, t, gt] = phase(w, n, id);
    G[v].push_back({u, int(G[u].size())-1, 0});
                                                                                                             id.erase(ranges::find(id, t));
                                                                                                             cut = min(cut, gt);
   Cap max_flow(int st_, int ed_) {
                                                                                                             for (int j = 0; j < n; ++j)</pre>
    st = st_, ed = ed_; Cap ret = 0;
for (int i = 63; i >= 0; --i)
                                                                                                               w[s][j] += w[t][j], w[j][s] += w[j][t];
      while (BFS(i)) ret += DFS(st);
                                                                                                            return cut;
     return ret;
                                                                                                         \frac{1}{V} = \frac{1}{V} \left( \frac{V^3}{V^3} \right), can be 0(VE + V^2 \log V)?
       // test @ luogu P3376
                                                                                                                     GomoryHu Tree [5edb29]
 4.5 HLPP [198e4e]
                                                                                                         vector<tuple<int, int, int>> GomoryHu(int n){
| template <typename T> struct HLPP {
                                                                                                          vector<tuple<int, int, int>> rt;
```

}

```
vector<int> g(n);
                                                             if (d[T] == INF_C) return nullopt;
for (int i = 1; i < n; ++i) {
                                                             for (size_t i = 0; i < d.size(); ++i) h[i] += d[i];</pre>
                                                             for (int i = T; i != S; i = f[i].first) {
 int t = g[i];
 auto f = flow;
                                                              auto &eg = g[f[i].first][f[i].second];
 rt.emplace_back(f.max_flow(i, t), i, t);
                                                              eg.f -= up[T]; g[eg.to][eg.r].f += up[T];
 f.walk(i); // bfs points that connected to i (use
    edges with .cap > 0)
                                                             return pair{up[T], h[T]};
 for (int j = i + 1; j < n; ++j)</pre>
   if (g[j]==t&&f.connect(j)) // check if i can reach j
                                                           public:
                                                            MCMF(int n) : g(n), f(n), up(n), d(n, INF_C) {}
   g[j] = i;
}
                                                            void add_edge(int s, int t, F c, C w) {
                                                             g[s].emplace_back(t, int(g[t].size()), c, w);
return rt;
                                                             g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
}
/* for our dinic:
* void walk(int) { BFS(0); }
                                                            pair<F, C> solve(int a, int b) {
* bool connect(int i) { return lv[i]; } */
                                                             h.assign(g.size(), 0);
                                                             F c = 0; C w = 0;
     Minimum Cost Max Flow [04f9cb]
                                                             while (auto r = step(a, b)) {
template <typename F, typename C> class MCMF {
                                                              c += r->first, w += r->first * r->second;
static constexpr F INF_F = numeric_limits<F>::max();
                                                              fill(d.begin(), d.end(), INF_C);
static constexpr C INF_C = numeric_limits<C>::max();
struct E { int to, r; F f; C c; };
                                                             return {c, w};
vector<vector<E>> g; vector<pair<int, int>> f;
vector<bool> inq; vector<F> up; vector<C> d;
                                                           };
optional<pair<F, C>> step(int S, int T) {
 queue<int> q;
                                                           4.10 Minimum Cost Circulation [3f7d84]
 for (q.push(S), d[S] = 0, up[S] = INF_F;
                                                           template <typename F, typename C>
   not q.empty(); q.pop()) {
                                                           struct MinCostCirculation {
  int u = q.front(); inq[u] = false;
                                                            struct ep { int to; F flow; C cost; };
  if (up[u] == 0) continue;
for (int i = 0; i < int(g[u].size()); ++i) {</pre>
                                                            int n; vector<int> vis; int visc;
                                                            vector<int> fa, fae; vector<vector<int>> g;
    auto e = g[u][i]; int v = e.to;
                                                            vector<ep> e; vector<C> pi;
    if (e.f <= 0 or d[v] <= d[u] + e.c) continue;</pre>
                                                            MinCostCirculation(int n_) : n(n_), vis(n), visc(0), g
   d[v] = d[u] + e.c; f[v] = \{u, i\};
                                                                (n), pi(n) {}
   up[v] = min(up[u], e.f);
                                                            void add_edge(int u, int v, F fl, C cs) {
   if (not inq[v]) q.push(v);
                                                             g[u].emplace_back((int)e.size());
    inq[v] = true;
                                                             e.emplace_back(v, fl, cs);
  }
                                                             g[v].emplace_back((int)e.size());
                                                             e.emplace_back(u, 0, -cs);
 if (d[T] == INF_C) return nullopt;
 for (int i = T; i != S; i = f[i].first) {
                                                            C phi(int x) {
  auto &eg = g[f[i].first][f[i].second];
                                                             if (fa[x] == -1) return 0;
  eg.f -= up[T]; g[eg.to][eg.r].f += up[T];
                                                             if (vis[x] == visc) return pi[x];
                                                             vis[x] = visc;
 return pair{up[T], d[T]};
                                                             return pi[x] = phi(fa[x]) - e[fae[x]].cost;
public:
                                                            int lca(int u, int v) {
MCMF(int n) : g(n),f(n),inq(n),up(n),d(n,INF_C) {}
                                                             for (; u != -1 || v != -1; swap(u, v)) if (u != -1) {
   if (vis[u] == visc) return u;
void add_edge(int s, int t, F c, C w) {
 g[s].emplace_back(t, int(g[t].size()), c, w);
                                                              vis[u] = visc;
 g[t].emplace_back(s, int(g[s].size()) - 1, 0, -w);
                                                              u = fa[u];
pair<F, C> solve(int a, int b) {
                                                             return -1;
 F c = 0; C w = 0;
 while (auto r = step(a, b)) {
                                                            void pushflow(int x, C &cost) {
  c += r->first, w += r->first * r->second;
                                                             int v = e[x ^ 1].to, u = e[x].to;
  ranges::fill(inq, false); ranges::fill(d, INF_C);
                                                              ++visc;
                                                             if (int w = lca(u, v); w == -1) {
 return {c, w};
                                                              while (v != -1)
}
                                                               swap(x ^= 1, fae[v]), swap(u, fa[v]), swap(u, v);
                                                             else {
      Dijkstra Cost Flow [d0cfd9]
                                                               int z = u, dir = 0; F f = e[x].flow;
template <typename F, typename C> class MCMF {
                                                              vector<int> cyc = {x};
static constexpr F INF_F = numeric_limits<F>::max();
                                                               for (int d : {0, 1})
static constexpr C INF_C = numeric_limits<C>::max();
                                                               for (int i = (d ? u : v); i != w; i = fa[i]) {
struct E { int to, r; F f; C c; };
                                                                cyc.push_back(fae[i] ^ d);
vector<vector<E>> g; vector<pair<int, int>> f;
                                                                if (chmin(f, e[fae[i] ^ d].flow)) z = i, dir = d;
vector<F> up; vector<C> d, h;
optional<pair<F, C>> step(int S, int T) {
                                                              for (int i : cyc) {
                                                               e[i].flow -= f; e[i ^ 1].flow += f;
 priority_queue<pair<C, int>> q;
 q.emplace(d[S] = 0, S), up[S] = INF_F;
                                                               cost += f * e[i].cost;
 while (not q.empty()) {
   auto [l, u] = q.top(); q.pop();
                                                              if (dir) x ^= 1, swap(u, v);
   if (up[u] == 0 or l != -d[u]) continue;
                                                              while (u != z)
   for (int i = 0; i < int(g[u].size()); ++i) {</pre>
                                                               swap(x ^= 1, fae[v]), swap(u, fa[v]), swap(u, v);
   auto e = g[u][i]; int v = e.to;
    auto nd = d[u] + e.c + h[u] - h[v];
    if (e.f <= 0 or d[v] <= nd) continue;</pre>
                                                            void dfs(int u) {
    f[v] = \{u, i\}; up[v] = min(up[u], e.f);
                                                             vis[u] = visc;
    q.emplace(-(d[v] = nd), v);
                                                             for (int i : g[u])
```

if (int v = e[i].to; vis[v] != visc and e[i].flow)

fa[v] = u, fae[v] = i, dfs(v);

void set_slack(int x) {

```
slack[x] = 0;
 C simplex() {
                                                                for (int u = 1; u <= n; ++u)</pre>
                                                                 if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
  C cost = 0;
  fa.assign(g.size(), -1); fae.assign(e.size(), -1);
                                                                  update_slack(u, x, slack[x]);
  ++visc; dfs(0);
  for (int fail = 0; fail < ssize(e); )</pre>
                                                               void q_push(int x) {
   for (int i = 0; i < ssize(e); i++)</pre>
                                                                if (x <= n) q.push(x);
                                                                else for (int y : flo[x]) q_push(y);
    if (e[i].flow and e[i].cost < phi(e[i ^ 1].to) -</pre>
    phi(e[i].to))
                                                               void set_st(int x, int b) {
     fail = 0, pushflow(i, cost), ++visc;
    else ++fail;
                                                                st[x] = b;
                                                                if (x > n) for (int y : flo[x]) set_st(y, b);
  return cost;
};
                                                               vector<int> split_flo(auto &f, int xr) {
                                                                auto it = find(all(f), xr);
4.11
      General Graph Matching [5f2293]
                                                                if (auto pr = it - f.begin(); pr % 2 == 1)
struct Matching {
                                                                 reverse(1 + all(f)), it = f.end() - pr;
 queue<int> q; int ans, n;
                                                                auto res = vector(f.begin(), it);
 vector<int> fa, s, v, pre, match;
                                                                return f.erase(f.begin(), it), res;
 int Find(int u) {
  return u == fa[u] ? u : fa[u] = Find(fa[u]); }
                                                               void set_match(int u, int v) {
 int LCA(int x, int y) {
  static int tk = 0; tk++; x = Find(x); y = Find(y);
                                                                match[u] = g[u][v].v;
                                                                if (u <= n) return;</pre>
  for (;; swap(x, y)) if (x != n) {
                                                                int xr = flo_from[u][g[u][v].u];
   if (v[x] == tk) return x;
                                                                auto &f = flo[u], z = split_flo(f, xr);
   v[x] = tk;
                                                                rep(i, 0, int(z.size())-1) set_match(z[i], z[i ^ 1]);
   x = Find(pre[match[x]]);
                                                                set_match(xr, v); f.insert(f.end(), all(z));
  }
                                                               void augment(int u, int v) {
 void Blossom(int x, int y, int l) {
   for (; Find(x) != l; x = pre[y]) {
                                                                for (;;) {
                                                                 int xnv = st[match[u]]; set_match(u, v);
   pre[x] = y, y = match[x];
                                                                 if (!xnv) return;
   if (s[y] == 1) q.push(y), s[y] = 0;
                                                                 set_match(xnv, st[pa[xnv]]);
   for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
                                                                 u = st[pa[xnv]], v = xnv;
 }
 bool Bfs(auto &&g, int r) {
                                                               int lca(int u, int v) {
  iota(all(fa), 0); ranges::fill(s, -1);
                                                                static int t = 0; ++t;
  q = queue<int>(); q.push(r); s[r] = 0;
                                                                for (++t; u || v; swap(u, v)) if (u) {
  for (; !q.empty(); q.pop()) {
                                                                 if (vis[u] == t) return u;
   for (int x = q.front(); int u : g[x])
                                                                 vis[u] = t; u = st[match[u]];
    if (s[u] == -1) {
                                                                 if (u) u = st[pa[u]];
     if (pre[u] = x, s[u] = 1, match[u] == n) {
      for (int a = u, b = x, last;
                                                                return 0;
        b != n; a = last, b = pre[a])
       last = match[b], match[b] = a, match[a] = b;
                                                               void add_blossom(int u, int o, int v) {
      return true;
                                                                int b = int(find(n + 1 + all(st), 0) - begin(st));
                                                                lab[b] = 0, S[b] = 0; match[b] = match[o];
     q.push(match[u]); s[match[u]] = 0;
                                                                vector<int> f = {o};
    } else if (!s[u] && Find(u) != Find(x)) {
                                                                for (int x = u, y; x != o; x = st[pa[y]])
     int l = LCA(u, x);
                                                                 f.pb(x), f.pb(y = st[match[x]]), q_push(y);
     Blossom(x, u, l); Blossom(u, x, l);
                                                                reverse(1 + all(f));
                                                                for (int x = v, y; x != o; x = st[pa[y]])
                                                                f.pb(x), f.pb(y = st[match[x]]), q_push(y);
flo[b] = f; set_st(b, b);
  return false;
                                                                for (int x = 1; x <= nx; ++x)
 Matching(auto &&g) : ans(0), n(int(g.size())),
                                                                 g[b][x].w = g[x][b].w = 0;
 fa(n+1), s(n+1), v(n+1), pre(n+1, n), match(n+1, n) {
  for (int x = 0; x < n; ++x)</pre>
                                                                for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;</pre>
                                                                for (int xs : flo[b]) {
   if (match[x] == n) ans += Bfs(g, x);
                                                                 for (int x = 1; x <= nx; ++x)</pre>
 } // match[x] == n means not matched
                                                                  if (g[b][x].w == 0 || ED(g[xs][x]) < ED(g[b][x]))
}; // test @ yosupo judge
                                                                   g[b][x] = g[xs][x], g[x][b] = g[x][xs];
4.12 Weighted Matching [94ca35]
                                                                 for (int x = 1; x \le n; ++x)
                                                                  if (flo_from[xs][x]) flo_from[b][x] = xs;
#define pb emplace_back
#define rep(i, l, r) for (int i=(l); i<=(r); ++i)
struct WeightGraph { // 1-based
                                                                set_slack(b);
 static const int inf = INT_MAX;
 struct edge { int u, v, w; }; int n, nx;
                                                               void expand_blossom(int b) {
 vector<int> lab; vector<vector<edge>> g;
                                                                for (int x : flo[b]) set_st(x, x);
vector<int> slack, match, st, pa, S, vis;
vector<vector<int>> flo, flo_from; queue<int> q;
                                                                int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
                                                                for (int x : split_flo(flo[b], xr)) {
 WeightGraph(int n_{-}) : n(n_{-}), nx(n * 2), lab(nx + 1),
                                                                 if (xs == -1) { xs = x; continue; }
  g(nx + 1, vector < edge > (nx + 1)), slack(nx + 1),
                                                                 pa[xs] = g[x][xs].u; S[xs] = 1, S[x] = 0;
  flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
                                                                 slack[xs] = 0; set_slack(x); q_push(x); xs = -1;
  match = st = pa = S = vis = slack;
  rep(u, 1, n) rep(v, 1, n) g[u][v] = \{u, v, 0\};
                                                                for (int x : flo[b])
                                                                 if (x == xr) S[x] = 1, pa[x] = pa[b];
                                                                 else S[x] = -1, set_slack(x);
 int ED(edge e) {
 return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
                                                                st[b] = 0;
 void update_slack(int u, int x, int &s) {
  if (!s || ED(g[u][x]) < ED(g[s][x])) s = u; }</pre>
                                                               bool on_found_edge(const edge &e) {
```

if (int u = st[e.u], v = st[e.v]; S[v] == -1) {

```
int nu = st[match[v]]; pa[v] = e.u; S[v] = 1;
   slack[v] = slack[nu] = 0; S[nu] = 0; q_push(nu);
  } else if (S[v] == 0) {
   if (int o = lca(u, v)) add_blossom(u, o, v);
   else return augment(u, v), augment(v, u), true;
  return false;
bool matching() {
 ranges::fill(S, -1); ranges::fill(slack, 0);
  q = queue<int>();
  for (int x = 1; x <= nx; ++x)</pre>
  if (st[x] == x && !match[x])
   pa[x] = 0, S[x] = 0, q_push(x);
  if (q.empty()) return false;
  for (;;) {
   while (q.size()) {
    int u = q.front(); q.pop();
    if (S[st[u]] == 1) continue;
    for (int v = 1; v <= n; ++v)</pre>
     if (g[u][v].w > 0 && st[u] != st[v]) {
      if (ED(g[u][v]) != 0)
      update_slack(u, st[v], slack[st[v]]);
      else if (on_found_edge(g[u][v])) return true;
   int d = inf;
   for (int b = n + 1; b <= nx; ++b)</pre>
   if (st[b] == b && S[b] == 1)
    d = min(d, lab[b] / 2);
   for (int x = 1; x \le nx; ++x)
    if (int s = slack[x]; st[x] == x && s && S[x] <= 0)</pre>
     d = min(d, ED(g[s][x]) / (S[x] + 2));
   for (int u = 1; u <= n; ++u)
    if (S[st[u]] == 1) lab[u] += d;
    else if (S[st[u]] == 0) {
     if (lab[u] <= d) return false;</pre>
     lab[u] -= d;
   rep(b, n + 1, nx) if (st[b] == b && S[b] >= 0)
   lab[b] += d * (2 - 4 * S[b]);
   for (int x = 1; x <= nx; ++x)
    if (int s = slack[x]; st[x] == x &&
      s \& st[s] != x \& ED(g[s][x]) == 0)
     if (on_found_edge(g[s][x])) return true;
   for (int b = n + 1; b <= nx; ++b)
    if (st[b] == b && S[b] == 1 && lab[b] == 0)
     expand_blossom(b);
  return false;
}
pair<lld, int> solve() {
 ranges::fill(match, 0);
  rep(u, 0, n) st[u] = u, flo[u].clear();
  int w_max = 0;
  rep(u, 1, n) rep(v, 1, n) {
  flo_from[u][v] = (u == v ? u : 0);
  w_max = max(w_max, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
 int n_matches = 0; lld tot_weight = 0;
while (matching()) ++n_matches;
  rep(u, 1, n) if (match[u] \&\& match[u] < u)
  tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
void set_edge(int u, int v, int w) {
 g[u][v].w = g[v][u].w = w; }
     Math
5.1 Common Bounds
   |234567892050100 n |1001e31e61e91e121e15
\binom{2n}{n} 2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 le7 4e7 1.5e8
```

5.2 Equations

Stirling Number of the First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

```
• S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)
• S_1(n,i) = [x^i] \left(\prod_{i=0}^{n-1} (x+i)\right), use D&Q and taylor shift.
• S_1(i,k) = \frac{i!}{k!} \left[ x^i \right] \left( \sum_{j \ge 1} \frac{x^j}{j} \right)^k
```

Stirling Number of the Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into knonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

$$S_2(i,k) = \frac{i}{k!} [x_1^i] (e^x - 1)^k$$

Derivatives/Integrals

Integration by parts:
$$\int_a^b f(x)g(x)dx = [F(x)g(x)]_a^b - \int_a^b F(x)g'(x)dx$$
 $\frac{d}{dx}\sin^{-1}x = \frac{1}{\sqrt{1-x^2}}\frac{d}{dx}\cos^{-1}x = -\frac{1}{\sqrt{1-x^2}}\frac{d}{dx}\tan^{-1}x = \frac{1}{1+x^2}$ $\frac{d}{dx}\tan x = 1 + \tan^2x \int \tan ax = -\frac{\ln|\cos ax|}{a}$ $\int e^{-x^2} = \frac{\sqrt{\pi}}{2} \mathrm{erf}(x) \int x e^{ax} dx = \frac{e^{ax}}{a^2}(ax-1)$

$$a^b \equiv \begin{cases} a^{(b \mod \varphi(m)) + \varphi(m)} & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m$$

Pentagonal Number Theorem

$$\prod_{n=1}^{\infty} (1 - x^n) = \sum_{k=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2} = (\sum p(n)x^n)^{-1}$$

5.3 Extended FloorSum

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1)) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

5.4 Integer Division [cd017d]

if (b < 0) {

```
lld fdiv(lld a, lld b)
{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
lld cdiv(lld a, lld b)
{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
5.5 FloorSum [fb5917]
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
// @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
 llu ans = 0:
 while (true) {
  if (a >= m) ans += n*(n-1)/2 * (a/m), a %= m;
  if (b >= m) ans += n * (b/m), b %= m;
  if (llu y_max = a * n + b; y_max >= m) {
   n = (llu)(y_max / m), b = (llu)(y_max % m);
   swap(m, a);
  } else break;
 }
 return ans;
lld floor_sum(lld n, lld m, lld a, lld b) {
 llu ans = 0;
 if (a < 0) {
  llu a2 = (a \% m + m), d = (a2 - a) / m;
  ans -= 1ULL * n * (n - 1) / 2 * d; a = a2;
```

llu b2 = (b % m + m), d = (b2 - b) / m;

return ans + floor_sum_unsigned(n, m, a, b);

ans -= 1ULL * n * d; b = b2;

```
5.6 ModMin [253e4d]
                                                             while (check(w = (z * z - n + P) % P) != P - 1)
// min{k | l <= ((ak) mod m) <= r}
                                                              z = rnd() \% P;
                                                             const auto M = [P, w](auto &u, auto &v) {
optional<lu> mod_min(u32 a, u32 m, u32 l, u32 r) {
                                                              auto [a, b] = u; auto [c, d] = v;
return make_pair((a * c + b * d % P * w) % P,
if (a == 0) return l ? nullopt : 0;
if (auto k = llu(l + a - 1) / a; k * a <= r)
  return k;
                                                                 (a * d + b * c) % P);
 auto b = m / a, c = m % a;
if (auto y = mod_min(c, a, a - r % a, a - l % a))
                                                             pair<lld, lld> r(1, 0), e(z, 1);
 return (l + *y * c + a - 1) / a + *y * b;
                                                             for (int q = (P + 1) / 2; q; q >>= 1, e = M(e, e))
                                                              if (q & 1) r = M(r, e);
return nullopt;
                                                             return int(r.first); // sqrt(n) mod P where P is prime
5.7 Floor Monoid Product [416e89]
/* template <typename T>
                                                            5.12 FWT [f82550]
T brute(llu a, llu b, llu c, llu n, T U, T R) {
                                                            /* or convolution:
                                                             * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
for (llu i = 1, l = 0; i <= n; i++, res = res * R)
                                                             * and convolution:
 for (llu \ r = (a*i+b)/c; \ l < r; ++l) \ res = res * U;
                                                            * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
void fwt(int x[], int N, bool inv = false) {
7 */
                                                             for (int d = 1; d < N; d <<= 1)</pre>
template <typename T>
                                                              for (int s = 0; s < N; s += d * 2)
T euclid(llu a, llu b, llu c, llu n, T U, T R) {
                                                               for (int i = s; i < s + d; i++) {</pre>
if (!n) return T{};
                                                                int j = i + d, ta = x[i], tb = x[j];
                                                                x[i] = add(ta, tb);
if (b >= c)
  return mpow(U, b / c) * euclid(a, b % c, c, n, U, R);
                                                                x[j] = sub(ta, tb);
if (a >= c)
  return euclid(a % c, b, c, n, U, mpow(U, a / c) * R);
                                                             if (inv) {
llu m = (u128(a) * n + b) / c;
                                                              const int invn = modinv(N);
if (!m) return mpow(R, n);
                                                              for (int i = 0; i < N; i++)
return mpow(R, (c - b - 1) / a) * U
                                                               x[i] = mul(x[i], invn);
 * euclid(c, (c - b - 1) % a, a, m - 1, R, U)
  * mpow(R, n - (u128(c) * m - b - 1) / a);
                                                            }
                                                            5.13 Packed FFT [7c64ba]
// time complexity is O(log max(a, b, c))
                                                            int round2k(size_t n) {
// UUUU R UUUUU R ... UUU R 共 N 個 R,最後一個必是 R
                                                             int sz = 1; while (sz < int(n)) sz *= 2; return sz; }</pre>
// 一直到第 k 個 R 前總共有 (ak+b)/c 個 U
                                                            VL convolution(const VI &a, const VI &b) {
5.8 ax+by=gcd [d0cbdd]
                                                             const int sz = round2k(a.size() + b.size() - 1);
// ax+ny = 1, ax+ny == ax == 1 \ (mod \ n)
                                                              // Should be able to handle N <= 10^5, C <= 10^4
void exgcd(lld x, lld y, lld &g, lld &a, lld &b) {
  if (y == 0) g = x, a = 1, b = 0;
                                                             vector<P> v(sz);
                                                             for (size_t i = 0; i < a.size(); ++i) v[i].RE(a[i]);</pre>
else exgcd(y, x \% y, g, b, a), b = (x / y) * a;
                                                             for (size_t i = 0; i < b.size(); ++i) v[i].IM(b[i]);</pre>
                                                             fft(v.data(), sz, /*inv=*/false);
5.9 Chinese Remainder [d69e74]
                                                             auto rev = v; reverse(1 + all(rev));
// please ensure r_i\in[0,m_i)
                                                             for (int i = 0; i < sz; ++i) +
bool crt(lld &m1, lld &r1, lld m2, lld r2) {
                                                              P A = (v[i] + conj(rev[i])) / P(2, 0);
                                                              P B = (v[i] - conj(rev[i])) / P(0, 2);
 if (m2 > m1) swap(m1, m2), swap(r1, r2);
 lld g, a, b; exgcd(m1, m2, g, a, b);
                                                              v[i] = A * B;
 if ((r2 - r1) % g != 0) return false;
 m2 /= g; lld D = (r2 - r1) / g % m2 * a % m2;
                                                             VL c(sz); fft(v.data(), sz, /*inv=*/true);
 r1 += (D < 0 ? D + m2 : D) * m1; m1 *= m2;
                                                             for (int i = 0; i < sz; ++i) c[i] = roundl(RE(v[i]));</pre>
 assert (r1 >= 0 && r1 < m1);
                                                             return c;
  return true;
                                                            VI convolution_mod(const VI &a, const VI &b) {
5.10 DiscreteLog [86e463]
                                                             const int sz = round2k(a.size() + b.size() - 1);
                                                             vector<P> fa(sz), fb(sz);
template<typename Int>
                                                             for (size_t i = 0; i < a.size(); ++i)</pre>
Int BSGS(Int x, Int y, Int M) {
// x^? \equiv y (mod M)
                                                              fa[i] = P(a[i] & ((1 << 15) - 1), a[i] >> 15);
Int t = 1, c = 0, g = 1;
                                                             for (size_t i = 0; i < b.size(); ++i)</pre>
for (Int M_ = M; M_ > 0; M_ >>= 1) g = g * x % M;
                                                              fb[i] = P(b[i] & ((1 << 15) - 1), b[i] >> 15);
for (g = gcd(g, M); t % g != 0; ++c) {
                                                             fft(fa.data(), sz); fft(fb.data(), sz);
                                                             auto rfa = fa; reverse(1 + all(rfa));
 if (t == y) return c;
 t = t * x % M;
                                                             for (int i = 0; i < sz; ++i) fa[i] *= fb[i];</pre>
                                                             for (int i = 0; i < sz; ++i) fb[i] *= conj(rfa[i]);</pre>
                                                             fft(fa.data(), sz, true); fft(fb.data(), sz, true);
if (y % g != 0) return -1;
t /= g, y /= g, M /= g;
                                                             vector<int> res(sz);
                                                             for (int i = 0; i < sz; ++i) {</pre>
Int h = 0, gs = 1;
                                                              lld A = (lld)roundl(RE((fa[i] + fb[i]) / P(2, 0)));
for (; h * h < M; ++h) gs = gs * x % M;
                                                              lld C = (lld) roundl(IM((fa[i] - fb[i]) / P(0, 2)));
unordered_map<Int, Int> bs;
                                                              lld B = (lld)roundl(IM(fa[i])); B %= p; C %= p;
for (Int s = 0; s < h; bs[y] = ++s) y = y * x % M;
                                                              res[i] = (A + (B << 15) + (C << 30)) % p;
for (Int s = 0; s < M; s += h) {
 t = t * gs % M;
                                                             return res;
 if (bs.count(t)) return c + s + h - bs[t];
                                                            } // test @ yosupo judge with long double
                                                            5.14 CRT for arbitrary mod [e4dde7]
return -1;
                                                            const int mod = 1000000007;
5.11 Quadratic Residue [f0baec]
                                                            const int M1 = 985661441; // G = 3 for M1, M2, M3
int get_root(int n, int P) { // ensure 0 <= n < p
if (P == 2 or n == 0) return n;</pre>
                                                            const int M2 = 998244353;
                                                            const int M3 = 1004535809;
auto check = [&](lld x) {
                                                            int superBigCRT(lld A, lld B, lld C) {
 return modpow(int(x), (P - 1) / 2, P); };
                                                              static_assert (M1 < M2 && M2 < M3);</pre>
if (check(n) != 1) return -1;
                                                              constexpr lld r12 = modpow(M1, M2-2, M2);
mt19937 \text{ rnd}(7122); lld z = 1, w;
                                                              constexpr lld r13 = modpow(M1, M3-2, M3);
```

```
constexpr lld r23 = modpow(M2, M3-2, M3);
                                                               auto B = Sx(Mul(Dx(A), Inv(A), bit_ceil(A.size()*2)));
  constexpr lld M1M2 = 1LL * M1 * M2 % mod;
                                                               return B.resize(A.size()), B;
  B = (B - A + M2) * r12 % M2;
  C = (C - A + M3) * r13 % M3;
                                                              S Exp(const S &v) { // coef[0] == 0; res[0] == 1
 C = (C - B + M3) * r23 % M3;
                                                               return Newton(v, 1,
  return (A + B * M1 + C * M1M2) % mod;
                                                                [](S &X, S &A, int sz) {
                                                                 auto Y = X; Y.resize(sz / 2); Y = Ln(Y);
                                                                 fi(0, Y.size()) Y[i] = sub(A[i], Y[i]);
5.15 NTT / FFT [41c1f2]
                                                                 Y[0] = add(Y[0], 1); X = Mul(X, Y, sz); );
template <int mod, int G, int maxn> struct NTT {
 static_assert (maxn == (maxn & -maxn));
                                                              S Pow(S a, lld M) { // period mod*(mod-1)
 int roots[maxn];
                                                               assert(!a.empty() && a[0] != 0);
 NTT () {
                                                               const auto imul = [&a](int s) {
  int r = modpow(G, (mod - 1) / maxn);
                                                                for (int &x: a) x = mul(x, s); }; int c = a[0];
  for (int i = maxn >> 1; i; i >>= 1) {
                                                               imul(modinv(c)); a = Ln(a); imul(int(M % mod));
   roots[i] = 1;
                                                               a = Exp(a); imul(modpow(c, int(M % (mod - 1))));
   for (int j = 1; j < i; j++)
                                                               return a; // mod x^N where N=a.size()
    roots[i + j] = mul(roots[i + j - 1], r);
   r = mul(r, r);
                                                              S Sqrt(const S &v) { // need: QuadraticResidue
   // for (int j = 0; j < i; j++) // FFT (tested)
                                                               assert(!v.empty() && v[0] != 0);
     roots[i+j] = polar < llf > (1, PI * j / i);
                                                               const int r = get_root(v[0]); assert(r != -1);
                                                               return Newton(v, r,
  [](S &X, S &A, int sz) {
 }
 // n must be 2^k, and 0 \le F[i] \le mod
                                                                 auto Y = X; Y.resize(sz / 2);
 template <typename T>
                                                                 auto B = Mul(A, Inv(Y), sz);
for (int i = 0, inv2 = mod / 2 + 1; i < sz; i++)</pre>
 void operator()(int F[], T n, bool inv = false) {
  for (T i = 0, j = 0; i < n; i++) {
                                                                  X[i] = mul(inv2, add(X[i], B[i])); });
   if (i < j) swap(F[i], F[j]);</pre>
   for (T k = n > 1; (j^k < k; k > 1);
                                                              S Mul(auto &&a, auto &&b) {
                                                               const auto n = a.size() + b.size() - 1;
  for (T s = 1; s < n; s *= 2) {
                                                               auto R = Mul(a, b, bit_ceil(n));
   for (T i = 0; i < n; i += s * 2) {
                                                               return R.resize(n), R;
    for (T j = 0; j < s; j++) {
  int a = F[i+j], b = mul(F[i+j+s], roots[s+j]);</pre>
                                                              S MulT(S a, S b, size_t k) {
     F[i+j] = add(a, b); // a + b
                                                               assert(b.size()); reverse(all(b)); auto R = Mul(a, b);
     F[i+j+s] = sub(a, b); // a - b
                                                               R = vector(R.begin() + b.size() - 1, R.end());
                                                               return R.resize(k), R;
   }
                                                              S Eval(const S &f, const S &x) {
  if (inv) {
                                                               if (f.empty()) return vector(x.size(), 0);
   int iv = modinv(int(n));
                                                               const int n = int(max(x.size(), f.size()));
   for (T i = 0; i < n; i++) F[i] = mul(F[i], iv);</pre>
                                                               auto q = vector(n * 2, S(2, 1)); S ans(n);
fi(0, x.size()) q[i + n][1] = sub(0, x[i]);
   reverse(F + 1, F + n);
                                                               for (int i = n - 1; i > 0; i--)
}
                                                                q[i] = Mul(q[i << 1], q[i << 1 | 1]);
};
                                                               q[1] = MulT(f, Inv(q[1]), n);
5.16 Formal Power Series [c6b99a]
                                                               for (int i = 1; i < n; i++) {</pre>
#define fi(l, r) for (size_t i = (\overline{l}); i < \overline{r}); i++)
                                                                auto L = q[i << 1], R = q[i << 1 | 1];</pre>
using S = vector<int>;
                                                                q[i << 1 | 0] = MulT(q[i], R, L.size());
                                                                q[i << 1 | 1] = MulT(q[i], L, R.size());</pre>
auto Mul(auto a, auto b, size_t sz) {
 a.resize(sz), b.resize(sz);
ntt(a.data(), sz); ntt(b.data(), sz);
fi(0, sz) a[i] = mul(a[i], b[i]);
                                                               for (int i = 0; i < n; i++) ans[i] = q[i + n][0];</pre>
                                                               return ans.resize(x.size()), ans;
 return ntt(a.data(), sz, true), a;
                                                              pair<S, S> DivMod(const S &A, const S &B) {
                                                               assert(!B.empty() && B.back() != 0);
S Newton(const S &v, int init, auto &&iter) {
 S Q = { init };
                                                               if (A.size() < B.size()) return {{}, A};</pre>
 for (int sz = 2; Q.size() < v.size(); sz *= 2) {</pre>
                                                               const auto sz = A.size() - B.size() + 1;
                                                               S X = B; reverse(all(X)); X.resize(sz);
  S A{begin(v), begin(v) + min(sz, int(v.size()))};
  A.resize(sz * 2), Q.resize(sz * 2);
                                                               S Y = A; reverse(all(Y)); Y.resize(sz);
                                                               S Q = Mul(Inv(X), Y);
  iter(Q, A, sz * 2); Q.resize(sz);
                                                               Q.resize(sz); reverse(all(Q)); X = Mul(Q, B); Y = A;
                                                               fi(0, Y.size()) Y[i] = sub(Y[i], X[i]);
return Q.resize(v.size()), Q;
                                                               while (Y.size() && Y.back() == 0) Y.pop_back();
S Inv(const S &v) { // v[0] != 0
                                                               while (Q.size() && Q.back() == 0) Q.pop_back();
return Newton(v, modinv(v[0]),
                                                               return {Q, Y};
                                                              } // empty means zero polynomial
  [](S &X, S &A, int sz) {
   ntt(X.data(), sz), ntt(A.data(), sz);
for (int i = 0; i < sz; i++)</pre>
                                                              int LinearRecursionKth(S a, S c, int64_t k) {
                                                               const auto d = a.size(); assert(c.size() == d + 1);
                                                               const auto sz = bit_ceil(2 * d + 1), o = sz / 2;
    X[i] = mul(X[i], sub(2, mul(X[i], A[i])));
                                                               S q = c; for (int &x: q) x = sub(0, x); q[0]=1;
   ntt(X.data(), sz, true); });
                                                               S p = Mul(a, q); p.resize(sz); q.resize(sz);
S Dx(S A) {
                                                               for (int r; r = (k & 1), k; k >>= 1) {
 fi(1, A.size()) A[i - 1] = mul(i, A[i]);
                                                                fill(d + all(p), 0); fill(d + 1 + all(q), 0);
 return A.empty() ? A : (A.pop_back(), A);
                                                                ntt(p.data(), sz); ntt(q.data(), sz);
                                                                for (size_t i = 0; i < sz; i++)</pre>
S Sx(S A) {
                                                                 p[i] = mul(p[i], q[(i + o) & (sz - 1)]);
                                                                for (size_t i = 0, j = 0; j < sz; i++, j++)</pre>
 A.insert(A.begin(), 0);
                                                                 q[i] = q[j] = mul(q[i], q[j]);
 fi(1, A.size()) A[i] = mul(modinv(int(i)), A[i]);
                                                                ntt(p.data(), sz, true); ntt(q.data(), sz, true);
 return A;
                                                                for (size_t i = 0; i < d; i++) p[i] = p[i << 1 | r];</pre>
                                                                for (size_t i = 0; i <= d; i++) q[i] = q[i << 1];</pre>
S Ln(const S &A) { // coef[0] == 1; res[0] == 0
```

```
} // Bostan-Mori
                                                                return t == n ? pollard_rho(n) : t;
                                                               } // test @ yosupo judge, ~270ms for Q=100
// if use montgomery, ~70ms for Q=100
 return mul(p[0], modinv(q[0]));
} // a_n = \sum_{j=0}^{n} a_{j-1}, c_0 \text{ is not used}
5.17 Partition Number [9bb845]
                                                               5.21 Berlekamp Massey [a94d00]
ans[0] = tmp[0] = 1;
for (int i = 1; i * i <= n; i++) {
                                                               template \langle typename T \rangle
                                                               vector<T> BerlekampMassey(const vector<T> &output) {
for (int rep = 0; rep < 2; rep++)
for (int j = i; j <= n - i * i; j++)</pre>
                                                                vector<T> d(output.size() + 1), me, he;
                                                                for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
 modadd(tmp[j], tmp[j-i]);
for (int j = i * i; j <= n; j++)</pre>
                                                                 for (size_t j = 0; j < me.size(); ++j)
d[i] += output[i - j - 2] * me[j];</pre>
 modadd(ans[j], tmp[j - i * i]);
                                                                 if ((d[i] -= output[i - 1]) == 0) continue;
                                                                 if (me.empty()) {
5.18
       Pi Count [715863]
                                                                  me.resize(f = i);
struct S { int rough; lld large; int id; };
                                                                  continue;
lld PrimeCount(lld n) { // n \sim 10^{13} \Rightarrow < 1s
                                                                 vector<T> o(i - f - 1);
 if (n <= 1) return 0;
                                                                 T k = -d[i] / d[f]; o.push_back(-k);
 const int v = static_cast<int>(sqrtl(n)); int pc = 0;
                                                                 for (T x : he) o.push_back(x * k);
 vector<int> smalls(v + 1), skip(v + 1); vector<S> z;
                                                                 if (o.size() < me.size()) o.resize(me.size());</pre>
 for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
                                                                 for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
 for (int i : views::iota(0, (v + 1) / 2))
                                                                 if (i-f+he.size() >= me.size()) he = me, f = i;
  z.emplace_back(2*i+1, (n / (2*i+1) + 1) / 2, i);
                                                                 me = o;
 for (int p = 3; p <= v; ++p)</pre>
                                                                }
  if (smalls[p] > smalls[p - 1]) {
                                                                return me;
  const int q = p * p; ++pc;
  if (1LL * q * q > n) break;
  skip[p] = 1;
                                                               5.22 Gauss Elimination [1f5f8c]
  for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
                                                               using VI = vector<int>;
  int ns = 0;
                                                               using VVI = vector<VI>;
  for (auto e : z) if (!skip[e.rough]) {
                                                               pair<VI, VVI> gauss(VVI A, VI b) {
  lld d = 1LL * e.rough * p;
                                                                const int N = (int)A.size(), M = (int)A[0].size();
   e.large += pc - (d <= v ? z[smalls[d] - pc].large :</pre>
                                                                vector<int> depv, free(M, true); int rk = 0;
   smalls[n / d]);
                                                                for (int i = 0; i < M; ++i) {</pre>
   e.id = ns; z[ns++] = e;
                                                                 int p = -1;
 }
                                                                 for (int j = rk; j < N; ++j)</pre>
  z.resize(ns);
                                                                  if (p == -1 || abs(A[j][i]) > abs(A[p][i]))
  for (int j = v / p; j >= p; --j) {
                                                                   p = j;
   int c = smalls[j] - pc, e = min(j * p + p, v + 1);
                                                                 if (p == -1 || A[p][i] == 0) continue;
   for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
                                                                 swap(A[p], A[rk]); swap(b[p], b[rk]);
                                                                 const int inv = modinv(A[rk][i]);
                                                                 for (int &x : A[rk]) x = mul(x, inv);
 lld ans = z[0].large; z.erase(z.begin());
                                                                 b[rk] = mul(b[rk], inv);
 for (auto &[rough, large, k] : z) {
  const lld m = n / rough; --k;
                                                                 for (int j = 0; j < N; ++j) if (j != rk) {</pre>
                                                                  int z = A[j][i];
  ans -= large - (pc + k);

for (auto [p, _, l] : z)

if (l >= k || p * p > m) break;
                                                                  for (int k = 0; k < M; ++k)
                                                                   A[j][k] = sub(A[j][k], mul(z, A[rk][k]));
                                                                  b[j] = sub(b[j], mul(z, b[rk]));
   else ans += smalls[m / p] - (pc + l);
                                                                 depv.push_back(i); free[i] = false; ++rk;
 return ans;
} // test @ yosupo library checker w/ n=1e11, 68ms
                                                                for (int i = rk; i < N; i++)</pre>
5.19 Miller Rabin [fbd812]
                                                                 if (b[i] != 0) return {{}}, {{}}}; // not consistent
bool isprime(llu x) {
                                                                VI x(M); VVI h;
                                                                for (int i = 0; i < rk; i++) x[depv[i]] = b[i];</pre>
 auto witn = [&](llu a, int t) {
  for (llu a2; t--; a = a2) {
                                                                for (int i = 0; i < M; i++) if (free[i]) {</pre>
                                                                 h.emplace_back(M); h.back()[i] = 1;
   a2 = mmul(a, a, x);
                                                                 for (int j = 0; j < rk; j++)</pre>
   if (a2 == 1 && a != 1 && a != x - 1) return true;
                                                                  h.back()[depv[j]] = sub(0, A[j][i]);
  }
  return a != 1;
                                                                return {x, h}; // solution = x + span(h[i])
 if (x <= 2 || ~x & 1) return x == 2;
 int t = countr_zero(x-1); llu odd = (x-1) >> t;
                                                               5.23 Charateristic Polynomial [cd559d]
 for (llu m:
                                                               #define rep(x, y, z) for (int x=y; x<z; x++)
  {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
                                                               using VI = vector<int>; using VVI = vector<VI>;
  if (m % x != 0 && witn(mpow(m % x, odd, x), t))
                                                               void Hessenberg(VVI &H, int N) {
   return false;
                                                                for (int i = 0; i < N - 2; ++i) {
 return true;
                                                                 for (int j = i + 1; j < N; ++j) if (H[j][i]) {</pre>
} // test @ luogu 143 & yosupo judge, ~1700ms for Q=1e5
                                                                  rep(k, i, N) swap(H[i+1][k], H[j][k]);
// if use montgomery, ~250ms for Q=1e5
                                                                  rep(k, 0, N) swap(H[k][i+1], H[k][j]);
5.20 Pollard Rho [57ad88]
                                                                  break:
// does not work when n is prime or n == 1
// return any non-trivial factor
                                                                 if (!H[i + 1][i]) continue;
llu pollard_rho(llu n) {
                                                                 for (int j = i + 2; j < N; ++j) {</pre>
 static mt19937_64 rnd(120821011);
                                                                  int co = mul(modinv(H[i + 1][i]), H[j][i]);
 if (!(n & 1)) return 2;
                                                                  rep(k, i, N) subeq(H[j][k], mul(H[i+1][k], co));
 llu y = 2, z = y, c = rnd() % n, p = 1, i = 0, t;
                                                                  rep(k, 0, N) addeq(H[k][i+1], mul(H[k][j], co));
 auto f = [&](llu x) {
  return madd(mmul(x, x, n), c, n); };
                                                                }
 do {
  p = mmul(msub(z = f(f(z)), y = f(y), n), p, n);
                                                               VI CharacteristicPoly(VVI A) {
  if (++i &= 63) if (i == (i & -i)) t = gcd(p, n);
                                                                int N = (int)A.size(); Hessenberg(A, N);
} while (t == 1);
                                                                VVI P(N + 1, VI(N + 1)); P[0][0] = 1;
```

```
for (int i = 1; i <= N; ++i) {</pre>
  rep(j, 0, i+1) P[i][j] = j ? P[i-1][j-1] : 0;
  for (int j = i - 1, val = 1; j >= 0; --j) {
   int co = mul(val, A[j][i - 1]);
   rep(k, 0, j+1) subeq(P[i][k], mul(P[j][k], co));
   if (j) val = mul(val, A[j][j - 1]);
 if (N \& 1) for (int \&x: P[N]) x = sub(0, x);
 return P[N]; // test: 2021 PTZ Korea K
5.24 Simplex [c9c93b]
namespace simplex {
// maximize c^Tx under Ax \le B and x \ge 0
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
using VD = vector<llf>;
using VVD = vector<vector<llf>>;
const llf eps = 1e-9, inf = 1e+9;
int n, m; VVD d; vector<int> p, q;
void pivot(int r, int s) {
 llf inv = 1.0 / d[r][s];
 for (int i = 0; i < m + 2; ++i)
  for (int j = 0; j < n + 2; ++j)</pre>
   if (i != r && j != s)
    d[i][j] -= d[r][j] * d[i][s] * inv;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;</pre>
 for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
 d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
 int x = m + z;
                                                                6
 while (true) {
  int s = -1;
  for (int i = 0; i <= n; ++i) {</pre>
   if (!z && q[i] == -1) continue;
   if (s == -1 || d[x][i] < d[x][s]) s = i;
  if (s == -1 || d[x][s] > -eps) return true;
  int r = -1;
  for (int i = 0; i < m; ++i) {</pre>
   if (d[i][s] < eps) continue;</pre>
   if (r == -1 ||
    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
  pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
m = (int)b.size(), n = (int)c.size();
 d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)</pre>
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)</pre>
  p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
 for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];
 q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
 for (int i = 1; i < m; ++i)</pre>
  if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
 if (d[r][n + 1] < -eps) {</pre>
  pivot(r, n);
  if (!phase(1) || d[m + 1][n + 1] < -eps)
  return VD(n, -inf);
for (int i = 0; i < m; ++i) if (p[i] == -1) {</pre>
   int s = min_element(d[i].begin(), d[i].end() - 1)
        - d[i].begin();
   pivot(i, s);
 }
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)</pre>
  if (p[i] < n) x[p[i]] = d[i][n + 1];</pre>
}} // use double instead of long double if possible
```

5.25 Simplex Construction

Standard form: maximize $\sum_{1 \le i \le n} c_i x_i$ such that $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$ for all $1 \le j \le m$ and $x_i \ge 0$ for all $1 \le i \le n$.

```
1. In case of minimization, let c_i' = -c_i
2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
3. \sum_{1 \leq i \leq n} A_{ji} x_i = b_j \rightarrow \mathsf{add} \leq \mathsf{and} \geq.
4. If x_i has no lower bound, replace x_i with x_i - x_i'
5.26 Adaptive Simpson [09669e]
llf simp(llf l, llf r) {
 llf m = (l + r) / 2;
 return (f(l) + f(r) + 4.0 * f(m)) * (r - l) / 6.0;
llf F(llf L, llf R, llf v, llf eps) {
 llf M = (L + R) / 2, vl = simp(L, M), vr = simp(M, R);
 if (abs(vl + vr - v) <= 15 * eps)
return vl + vr + (vl + vr - v) / 15.0;</pre>
 return F(L, M, vl, eps / 2.0) +
     F(M, R, vr, eps / 2.0);
     call F(l, r, simp(l, r), 1e-6)
5.27 Golden Ratio Search [376bcb]
llf gss(llf a, llf b, auto &&f) {
  llf r = (sqrt(5)-1)/2, eps = 1e-7;
llf x1 = b - r*(b-a), x2 = a + r*(b-a);
  llf f1 = f(x1), f2 = f(x2);
  while (b-a > eps)
    if (f1 < f2) { //change to > to find maximum
      b = x2; x2 = x1; f2 = f1;
      x1 = b - r*(b-a); f1 = f(x1);
    } else {
      a = x1; x1 = x2; f1 = f2;
      x2 = a + r*(b-a); f2 = f(x2);
    }
  return a;
     Geometry
     Basic Geometry [1d2d70]
#define IM imag
#define RE real
using lld = int64_t;
using llf = long double;
using PT = complex<lld>;
using PF = complex<llf>;
using P = PT;
llf abs(P p) { return sqrtl(norm(p)); }
PF toPF(PT p) { return PF{RE(p), IM(p)}; }
int sgn(lld x) { return (x > 0) - (x < 0); }</pre>
lld dot(P a, P b) { return RE(conj(a) * b); }
lld cross(P a, P b) { return IM(conj(a) * b); }
int ori(P a, P b, P c) {
 return sgn(cross(b - a, c - a));
int quad(P p) {
 return (IM(p) == 0) // use sgn for PF
  ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
int argCmp(P a, P b) {
 // returns 0/+-1, starts from theta = -PI
 int qa = quad(a), qb = quad(b);
 if (qa != qb) return sgn(qa - qb);
 return sgn(cross(b, a));
P rot90(P p) { return P{-IM(p), RE(p)}; }
template <typename V> llf area(const V & pt) {
 lld ret = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 return ret / 2.0;
template <typename V> PF center(const V & pt) {
 P ret = 0; lld A = 0;
 for (int i = 1; i + 1 < (int)pt.size(); i++) {</pre>
  lld cur = cross(pt[i] - pt[0], pt[i+1] - pt[0]);
  ret += (pt[i] + pt[i + 1] + pt[0]) * cur; A += cur;
 return toPF(ret) / llf(A * 3);
PF project(PF p, PF q) { // p onto q
 return dot(p, q) * q / dot(q, q); // dot<llf>
6.2 2D Convex Hull [ecba37]
// from NaCl, counterclockwise, be careful of n<=2
vector<P> convex_hull(vector<P> v) {
 sort(all(v)); // by X then Y
 if (v[0] == v.back()) return {v[0]};
```

```
int t = 0, s = 1; vector<P> h(v.size() + 1);
                                                                  return U::valid(cross(D, B.dir), C) &&
 for (int _ = 2; _-
                     -; s = t--, reverse(all(v)))
                                                                  V::valid(cross(D, A.dir), C);
  for (P p : v) {
   while (t>s && ori(p, h[t-1], h[t-2]) >= 0) t--;
                                                                 6.7 Half Plane Intersection [31e216]
   h[t++] = p;
                                                                struct Line {
  }
                                                                 P st, ed, dir;
 return h.resize(t), h;
                                                                 Line (P s, P e) : st(s), ed(e), dir(e - s) {}
                                                                }; using LN = const Line &;
      2D Farthest Pair [8b5844]
                                                                PF intersect(LN A, LN B) {
                                                                 llf t = cross(B.st - A.st, B.dir) /
// p is CCW convex hull w/o colinear points
int n = (int)p.size(), pos = 1; lld ans = 0;
                                                                  llf(cross(A.dir, B.dir));
for (int i = 0; i < n; i++) {</pre>
                                                                  return toPF(A.st) + toPF(A.dir) * t; // C^3 / C^2
P = p[(i + 1) \% n] - p[i];
 while (cross(e, p[(pos + 1) % n] - p[i]) >
                                                                bool cov(LN l, LN A, LN B) {
     cross(e, p[pos] - p[i]))
                                                                  i128 u = cross(B.st-A.st, B.dir);
  pos = (pos + 1) % n;
                                                                  i128 v = cross(A.dir, B.dir);
 for (int j: {i, (i + 1) % n})
ans = max(ans, norm(p[pos] - p[j]));
                                                                  // ori(l.st, l.ed, A.st + A.dir*(u/v)) <= 0?
                                                                 i128 x = RE(A.dir) * u + RE(A.st - l.st) * v;
i128 y = IM(A.dir) * u + IM(A.st - l.st) * v;
} // tested @ AOJ CGL_4_B
                                                                  return sgn(x*IM(l.dir) - y*RE(l.dir)) * sgn(v) >= 0;
6.4 MinMax Enclosing Rect [e4470c]
                                                                 } // x, y are C^3, also sgn<i128> is needed
// from 8BQube, plz ensure p is strict convex hull
const llf INF = 1e18, qi = acos(-1) / 2 * 3;
                                                                bool operator<(LN a, LN b) {</pre>
                                                                 if (int c = argCmp(a.dir, b.dir)) return c == -1;
pair<llf, llf> solve(const vector<P> &p) {
                                                                  return ori(a.st, a.ed, b.st) < 0;</pre>
 llf mx = 0, mn = INF; int n = (int)p.size();
 for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {
                                                                // cross(pt-line.st, line.dir)<=0 <-> pt in half plane
#define Z(v) (p[(v) % n] - p[i])
  P e = Z(i + 1);
                                                                 // the half plane is the LHS when going from st to ed
                                                                llf HPI(vector<Line> &q) {
  while (cross(e, Z(u + 1)) > cross(e, Z(u))) ++u;
  while (dot(e, Z(r + 1)) > dot(e, Z(r))) ++r;
if (!i) l = r + 1;
                                                                  sort(q.begin(), q.end());
                                                                 int n = (int)q.size(), l = 0, r = -1;
for (int i = 0; i < n; i++) {
   if (i && !argCmp(q[i].dir, q[i-1].dir)) continue;</pre>
  while (dot(e, Z(l + 1)) < dot(e, Z(l))) ++l;</pre>
 P D = p[r % n] - p[l % n];
llf H = cross(e, Z(u)) / llf(norm(e));
                                                                   while (l < r && cov(q[i], q[r-1], q[r])) --r;</pre>
                                                                  while (l < r && cov(q[i], q[l], q[l+1])) ++l;</pre>
  mn = min(mn, dot(e, D) * H);
  llf B = sqrt(norm(D)) * sqrt(norm(Z(u)));
                                                                  q[++r] = q[i];
  llf deg = (qi - acos(dot(D, Z(u)) / B)) / 2;
                                                                  while (l < r && cov(q[l], q[r-1], q[r])) --r;
  mx = max(mx, B * sin(deg) * sin(deg));
                                                                 while (l < r && cov(q[r], q[l], q[l+1])) ++l;
n = r - l + 1; // q[l .. r] are the lines</pre>
 return {mn, mx};
                                                                  if (n <= 1 || !argCmp(q[l].dir, q[r].dir)) return 0;</pre>
} // test @ UVA 819
                                                                  vector<PF> pt(n);
6.5 Minkowski Sum [602806]
                                                                  for (int i = 0; i < n; i++)</pre>
// A, B are strict convex hull rotate to min by (X, Y)
                                                                  pt[i] = intersect(q[i+l], q[(i+1)%n+l]);
vector<P> Minkowski(vector<P> A, vector<P> B) {
  const int N = (int)A.size(), M = (int)B.size();
                                                                  return area(pt);
                                                                } // test @ 2020 Nordic NCPC : BigBrother
vector<P> sa(N), sb(M), C(N + M + 1);
for (int i = 0; i < N; i++) sa[i] = A[(i+1)%N]-A[i];
for (int i = 0; i < M; i++) sb[i] = B[(i+1)%M]-B[i];</pre>
                                                                 6.8 SegmentDist (Sausage) [9d8603]
                                                                 // be careful of abs<complex<int>> (replace _abs below)
                                                                llf PointSegDist(P A, Seg B) {
 C[0] = A[0] + B[0];
 for (int i = 0, j = 0; i < N || j < M; ) {
  P e = (j>=M || (i<N && cross(sa[i], sb[j])>=0))
                                                                  if (B.dir == P(0)) return _abs(A - B.st);
                                                                  if (sgn(dot(A - B.st, B.dir)) *
                                                                    sgn(dot(A - B.ed, B.dir)) <= 0)</pre>
   ? sa[i++] : sb[j++];
                                                                   return abs(cross(A - B.st, B.dir)) / _abs(B.dir);
  C[i + j] = e;
                                                                  return min(_abs(A - B.st), _abs(A - B.ed));
 partial_sum(all(C), C.begin()); C.pop_back();
                                                                 llf SegSegDist(const Seg &s1, const Seg &s2) {
 return convex_hull(C); // just to remove colinear
                                                                  if (isInter(s1, s2)) return 0;
                                                                  return min({
      Segment Intersection [60d016]
6.6
                                                                    PointSegDist(s1.st, s2),
struct Seg { // closed segment
                                                                    PointSegDist(s1.ed, s2),
 P st, dir; // represent st + t*dir for 0<=t<=1
                                                                    PointSegDist(s2.st, s1),
 Seg(P s, P e) : st(s), dir(e - s) {}
                                                                    PointSegDist(s2.ed, s1) });
 static bool valid(lld p, lld q) {
                                                                } // test @ QOJ2444 / PTZ19 Summer.D3
  // is there t s.t. 0 <= t <= 1 && qt == p ?
                                                                 6.9 Rotating Sweep Line [8aff27]
  if (q < 0) q = -q, p = -p;
  return 0 <= p && p <= q;
                                                                struct Event {
                                                                 Pd; int u, v;
 vector<P> ends() const { return { st, st + dir }; }
                                                                  bool operator<(const Event &b) const {</pre>
                                                                   return sgn(cross(d, b.d)) > 0; }
template <typename T> bool isInter(T A, P p) {
 if (A.dir == P(0)) return p == A.st; // BE CAREFUL
                                                                P makePositive(P z) { return cmpxy(z, 0) ? -z : z; }
 return cross(p - A.st, A.dir) == 0 &&
                                                                 void rotatingSweepLine(const vector<P> &p) {
  T::valid(dot(p - A.st, A.dir), norm(A.dir));
                                                                  const int n = int(p.size());
                                                                  vector<Event> e; e.reserve(n * (n - 1) / 2);
                                                                  for (int i = 0; i < n; i++)
for (int j = i + 1; j < n; j++)</pre>
template <typename U, typename V>
bool isInter(U A, V B) {
 if (cross(A.dir, B.dir) == 0) { // BE CAREFUL
                                                                    e.emplace_back(makePositive(p[i] - p[j]), i, j);
 bool res = false
                                                                  sort(all(e));
  for (P p: A.ends()) res |= isInter(B, p);
                                                                  vector<int> ord(n), pos(n);
                                                                  iota(all(ord), 0);
  for (P p: B.ends()) res |= isInter(A, p);
                                                                  sort(all(ord), [&p](int i, int j) {
  return res;
                                                                  return cmpxy(p[i], p[j]); });
                                                                  for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
 P D = B.st - A.st; lld C = cross(A.dir, B.dir);
```

```
const auto makeReverse = [](auto &v) {
                                                              llf t = (x - RE(a.st)) / llf(RE(a.dir));
 sort(all(v)); v.erase(unique(all(v)), v.end());
                                                              return IM(a.st) + IM(a.dir) * t;
  vector<pair<int,int>> segs;
  for (size_t i = 0, j = 0; i < v.size(); i = j) {</pre>
                                                             lld cur_x = 0;
  for (; j < v.size() && v[j] - v[i] <= j - i; j++);</pre>
                                                             auto cmp = [&](const Seg &a, const Seg &b) -> bool {
  segs.emplace_back(v[i], v[j-1]+1+1);
                                                              if (int s = sgn(eval(a, cur_x) - eval(b, cur_x)))
                                                               return s == -1;
  return segs;
                                                              int s = sgn(cross(b.dir, a.dir));
                                                              if (cur_x != RE(a.st) && cur_x != RE(b.st)) s *= -1;
for (size_t i = 0, j = 0; i < e.size(); i = j) {</pre>
                                                              return s == -1;
 /* do here */
 vector<size_t> tmp;
                                                             namespace pbds = __gnu_pbds;
                                                             using Tree = pbds::tree<Seg, int, decltype(cmp),</pre>
 for (; j < e.size() && !(e[i] < e[j]); j++)</pre>
   tmp.push_back(min(pos[e[j].u], pos[e[j].v]));
                                                                pbds::rb_tree_tag,
 for (auto [l, r] : makeReverse(tmp)) {
                                                                pbds::tree_order_statistics_node_update>;
  reverse(ord.begin() + l, ord.begin() + r);
                                                             Tree st(cmp);
   for (int t = l; t < r; t++) pos[ord[t]] = t;</pre>
                                                             auto answer = [&](P ep) {
 }
                                                              if (binary_search(all(vtx), ep))
}
                                                               return 1; // on vertex
                                                              Seg H(ep, ep); // ??
                                                              auto it = st.lower_bound(H);
6.10
      Polygon Cut [fdd064]
                                                              if (it != st.end() && isInter(it->first, ep))
using P = PF;
                                                               return 1; // on edge
vector<P> cut(const vector<P>& poly, P s, P e) {
                                                              if (it != st.begin() && isInter(prev(it)->first, ep))
vector<P> res;
                                                               return 1; // on edge
for (size_t i = 0; i < poly.size(); i++) {</pre>
                                                              auto rk = st.order_of_key(H);
 P cur = poly[i], prv = i ? poly[i-1] : poly.back();
                                                              return rk % 2 == 0 ? 0 : 2; // 0: outside, 2: inside
 bool side = ori(s, e, cur) < 0;</pre>
  if (side != (ori(s, e, prv) < 0))
                                                             vector<int> ans(Q);
   res.push_back(intersect({s, e}, {cur, prv}));
                                                             for (auto [ep, i] : evt) {
 if (side)
                                                              cur_x = RE(ep);
  res.push_back(cur);
                                                              if (i < 0) { // remove
                                                               st.erase(edge[~i]);
return res;
                                                              } else if (i < N) { // insert</pre>
                                                               auto [it, succ] = st.insert({edge[i], i});
      Point In Simple Polygon [037c52]
                                                               assert (succ);
bool PIP(const vector<P> &p, P z, bool strict = true) {
                                                              } else
int cnt = 0, n = (int)p.size();
                                                               ans[i - N] = answer(ep);
for (int i = 0; i < n; i++) {</pre>
 P A = p[i], B = p[(i + 1) % n];
                                                             return ans;
  if (isInter(Seg(A, B), z)) return !strict;
                                                           } // test @ AOJ CGL_3_C
  auto zy = IM(z), Ay = IM(A), By = IM(B);
                                                            6.14 Tangent of Points To Hull [6d7cd7]
  cnt ^= ((zy<Ay) - (zy<By)) * ori(z, A, B) > 0;
                                                            pair<int, int> get_tangent(const vector<P> &v, P p) {
return cnt;
                                                             const auto gao = [&, N = int(v.size())](int s) {
                                                              const auto lt = [&](int x, int y) {
                                                              return ori(p, v[x % N], v[y % N]) == s; };
6.12 Point In Hull (Fast) [060ba1]
                                                              int l = 0, r = N; bool up = lt(0, 1);
bool PIH(const vector<P> &h, P z, bool strict = true) {
                                                              while (r - l > 1) {
int n = (int)h.size(), a = 1, b = n - 1, r = !strict;
                                                               int m = (l + r) / 2;
if (n < 3) return r && isInter(Seg(h[0], h[n-1]), z);</pre>
                                                               if (lt(m, 0) ? up : !lt(m, m+1)) r = m;
 if (ori(h[0],h[a],h[b]) > 0) swap(a, b);
                                                               else l = m:
if (ori(h[0],h[a],z) >= r || ori(h[0],h[b],z) <= -r)</pre>
  return false;
                                                              return (lt(l, r) ? r : l) % N;
while (abs(a - b) > 1) {
                                                             }; // test @ codeforces.com/gym/101201/problem/E
 int c = (a + b) / 2;
                                                             return {gao(-1), gao(1)}; // (a,b):ori(p,v[a],v[b])<0
  (ori(h[0], h[c], z) > 0 ? b : a) = c;
                                                            } // plz ensure that point strictly out of hull
return ori(h[a], h[b], z) < r;
                                                                  Circle Class & Intersection [d5df51]
                                                            llf FMOD(llf x) {
6.13
      Point In Polygon (Fast) [00590a]
                                                             if (x < -PI) x += PI * 2;
vector<int> PIPfast(vector<P> p, vector<P> q) {
                                                             if (x > PI) x -= PI * 2;
                                                             return x;
const int N = int(p.size()), Q = int(q.size());
vector<pair<P, int>> evt;
                                                            struct Cir { PF o; llf r; };
vector<Seg> edge;
 for (int i = 0; i < N; i++) {</pre>
                                                            // be carefule when tangent
 int a = i, b = (i + 1) % N;
                                                            vector<llf> intersectAngle(Cir a, Cir b) {
 P A = p[a], B = p[b];
                                                            PF dir = b.o - a.o; llf d2 = norm(dir);
 assert (A < B || B < A); // std::operator<
                                                             if (norm(a.r - b.r) >= d2) { // <math>norm(x) := |x|^2
 if (B < A) swap(A, B);
                                                             if (a.r < b.r) return {-PI, PI}; // a in b</pre>
 evt.emplace_back(A, i);
                                                              else return {}; // b in a
 evt.emplace_back(B, ~i);
                                                             } else if (norm(a.r + b.r) <= d2) return {};</pre>
                                                             llf dis = abs(dir), theta = arg(dir);
 edge.emplace_back(A, B);
                                                             llf phi = acos((a.r * a.r + d2 - b.r * b.r) /
for (int i = 0; i < Q; i++)</pre>
                                                               (2 * a.r * dis)); // is acos_safe needed ?
                                                            llf L = FMOD(theta - phi), R = FMOD(theta + phi);
 evt.emplace_back(q[i], i + N);
                                                             return { L, R };
 sort(all(evt));
auto vtx = p; sort(all(vtx));
auto eval = [](const Seg &a, lld x) -> llf {
                                                            vector<PF> intersectPoint(Cir a, Cir b) {
 if (RE(a.dir) == 0) {
                                                            llf d = abs(a.o - b.o);
  assert (x == RE(a.st));
                                                             if (d > b.r+a.r || d < abs(b.r-a.r)) return {};</pre>
   return IM(a.st) + llf(IM(a.dir)) / 2;
                                                             llf dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
                                                            PF dir = (a.o - b.o) / d;
```

```
PF u = dir * d1 + b.o;
                                                                    c.r = abs(p[i] - c.o);
                                                                    for (size_t k = 0; k < j; k++) {</pre>
 PF v = rot90(dir) * sqrt(max(0.0L, b.r*b.r-d1*d1));
                                                                     if (abs(p[k] - c.o) <= c.r) continue;</pre>
 return \{u + v, u - v\};
} // test @ AOJ CGL probs
                                                                     c = getCircum(p[i], p[j], p[k]);
6.16 Circle Common Tangent [d97f1c]
                                                                   }
// be careful of tangent / exact same circle
// sign1 = 1 for outer tang, -1 for inter tang
                                                                  return c;
vector<Line> common_tan(const Cir &a, const Cir &b, int
                                                                 } // test @ TIOJ 1093 & luogu P1742
     sign1) {
 if (norm(a.o - b.o) < eps) return {};</pre>
                                                                 6.20 Circle Union [073c1c]
 llf d = abs(a.o - b.o), c = (a.r - sign1 * b.r) / d;
                                                                 #define eb emplace_back
 PF v = (b.o - a.o) / d;
                                                                 struct Teve { // test@SPOJ N=1000, 0.3~0.5s
 if (c * c > 1) return {};
                                                                 PF p; llf a; int add; // point, ang, add
Teve(PF x, llf y, int z) : p(x), a(y), add(z) {}
 if (abs(c * c - 1) < eps) {
  PF p = a.o + c * v * a.r;
                                                                  bool operator<(Teve &b) const { return a < b.a; }</pre>
  return {Line(p, p + rot90(b.o - a.o))};
                                                                 // strict: x = 0, otherwise x = -1
 vector<Line> ret; llf h = sqrt(max(0.0L, 1-c*c));
for (int sign2 : {1, -1}) {
                                                                bool disjunct(Cir &a, Cir &b, int x)
                                                                { return sgn(abs(a.o - b.o) - a.r - b.r) > x; }
bool contain(Cir &a, Cir &b, int x)
  PF n = c * v + sign2 * h * rot90(v);
  PF p1 = a.o + n * a.r;
                                                                 { return sgn(a.r - b.r - abs(a.o - b.o)) > x; }
  PF p2 = b.o + n * (b.r * sign1);
                                                                 vector<llf> CircleUnion(vector<Cir> &c) {
  ret.emplace_back(p1, p2);
                                                                  // area[i] : area covered by at least i circles
 }
                                                                  int N = (int)c.size(); vector<llf> area(N + 1);
 return ret;
                                                                  vector<vector<int>> overlap(N, vector<int>(N));
                                                                  auto g = overlap; // use simple 2darray to speedup
6.17 Line-Circle Intersection [10786a]
                                                                  for (int i = 0; i < N; ++i)</pre>
vector<PF> LineCircleInter(PF p1, PF p2, PF o, llf r) {
                                                                   for (int j = 0; j < N; ++j) {
   /* c[j] is non-strictly in c[i]. */</pre>
 PF ft = p1 + project(o-p1, p2-p1), vec = p2-p1;
 llf dis = abs(o - ft);
                                                                    overlap[i][j] = i != j &&
 if (abs(dis - r) < eps) return {ft};</pre>
                                                                     (sgn(c[i].r - c[j].r) > 0 | | (sgn(c[i].r - c[j].r) == 0 && i < j)) &&
if (dis > r) return {};
vec = vec * sqrt(r * r - dis * dis) / abs(vec);
                                                                     contain(c[i], c[j], -1);
 return {ft + vec, ft - vec}; // sqrt_safe?
                                                                  for (int i = 0; i < N; ++i)</pre>
6.18 Poly-Circle Intersection [8e5133]
                                                                   for (int j = 0; j < N; ++j)</pre>
// Divides into multiple triangle, and sum up
                                                                    g[i][j] = i != j && !(overlap[i][j] ||
  from 8BQube, test by HDU2892 & AOJ CGL_7_H
                                                                  overlap[j][i] || disjunct(c[i], c[j], -1));
for (int i = 0; i < N; ++i) {</pre>
llf _area(PF pa, PF pb, llf r) {
 if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
                                                                   vector<Teve> eve; int cnt = 1;
 if (abs(pb) < eps) return 0;</pre>
                                                                   for (int j = 0; j < N; ++j) cnt += overlap[j][i];
// if (cnt > 1) continue; (if only need area[1])
 llf S, h, theta;
 llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
llf cB = dot(pb, pb-pa) / a / c, B = acos_safe(cB);
                                                                   for (int j = 0; j < N; ++j) if (g[i][j]) {</pre>
                                                                    auto IP = intersectPoint(c[i], c[j]);
 llf cC = dot(pa, pb) / a / b, C = acos_safe(cC);
                                                                    PF aa = IP[1], bb = IP[0];
 if (a > r) {
                                                                    llf A = arg(aa - c[i].o), B = arg(bb - c[i].o);
  S = (C / 2) * r * r; h = a * b * sin(C) / c;
                                                                    eve.eb(bb, B, 1); eve.eb(aa, A, -1);
  if (h < r && B < PI / 2)
                                                                    if (B > A) ++cnt;
   S = (acos\_safe(h/r)*r*r - h*sqrt\_safe(r*r-h*h));
 } else if (b > r) {
  theta = PI - B - asin_safe(sin(B) / r * a);
                                                                   if (eve.empty()) area[cnt] += PI*c[i].r*c[i].r;
                                                                   else {
  S = 0.5 * a*r*sin(theta) + (C-theta)/2 * r * r;
                                                                    sort(eve.begin(), eve.end());
                                                                    eve.eb(eve[0]); eve.back().a += PI * 2;
  S = 0.5 * sin(C) * a * b;
                                                                    for (size_t j = 0; j + 1 < eve.size(); j++) {</pre>
 return S;
                                                                     cnt += eve[j].add;
                                                                     area[cnt] += cross(eve[j].p, eve[j+1].p) *.5;
llf area_poly_circle(const vector<PF> &v, PF 0, llf r)
                                                                     llf t = eve[j + 1].a - eve[j].a;
                                                                     area[cnt] += (t-sin(t)) * c[i].r * c[i].r *.5;
 llf S = 0;
 for (size_t i = 0, N = v.size(); i < N; ++i)</pre>
                                                                   }
  S += _area(v[i] - 0, v[(i + 1) % N] - 0, r) *
    ori(0, v[i], v[(i + 1) % N]);
                                                                  return area;
 return abs(S);
6.19 Minimum Covering Circle [92bb15]
                                                                 6.21 Polygon Union [2bff43]
Cir getCircum(P a, P b, P c){ // P = complex<llf>
P z1 = a - b, z2 = a - c; llf D = cross(z1, z2) * 2;
                                                                llf rat(P \overline{a}, \overline{P} b) { return sgn(RE(b)) ? llf(RE(a))/RE(b
                                                                      ) : llf(IM(a))/IM(b); }
                                                                 llf polyUnion(vector<vector<P>>& poly) {
 auto c1 = dot(a + b, z1), c2 = dot(a + c, z2);
 P \circ = rot90(c2 * z1 - c1 * z2) / D;
                                                                  llf ret = 0; // area of poly[i] must be non-negative
 return { o, abs(o - a) };
                                                                  rep(i,0,sz(poly)) rep(v,0,sz(poly[i])) {
                                                                   P A = poly[i][v], B = poly[i][(v + 1) % sz(poly[i])];
Cir minCircleCover(vector<P> p) {
                                                                   vector<pair<llf, int>> segs{{0, 0}, {1, 0}};
 assert (!p.empty());
                                                                   rep(j,0,sz(poly)) if (i != j) {
 ranges::shuffle(p, mt19937(114514));
                                                                    rep(u,0,sz(poly[j])) {
                                                                     P C = poly[j][u], D = poly[j][(u + 1) % sz(poly[j])
 Cir c = { 0, 0 };
 for (size_t i = 0; i < p.size(); i++) {</pre>
                                                                     ];
if (int sc = ori(A, B, C), sd = ori(A, B, D); sc !=
  if (abs(p[i] - c.o) <= c.r) continue;</pre>
  c = { p[i], 0 };
                                                                      sd) {
  for (size_t j = 0; j < i; j++) {
   if (abs(p[j] - c.o) <= c.r) continue;</pre>
                                                                      llf sa = cross(D-C, A-C), sb = cross(D-C, B-C);
                                                                      if (min(sc, sd) < 0)
   c.o = (p[i] + p[j]) / llf(2);
                                                                       segs.emplace_back(sa / (sa - sb), sgn(sc - sd));
```

```
} else if (!sc && !sd && j<i && sgn(dot(B-A,D-C))</pre>
                                                                return now;
    >0){
     segs.emplace_back(rat(C - A, B - A), 1);
                                                               }
     segs.emplace_back(rat(D - A, B - A), -1);
                                                               // n^2 delaunay: facets with negative z normal of
                                                               // convexhull of (x, y, x^2 + y^2), use a pseudo-point
   }
                                                               // (0, 0, inf) to avoid degenerate case
                                                               // test @ SPOJ CH3D
                                                               // llf area = 0, vol = 0; // surface area / volume
  sort(segs.begin(), segs.end());
  for (auto &s : segs) s.first = clamp<llf>(s.first, 0,
                                                               // for (auto [a, b, c]: faces)
                                                                  area += abs(ver(p[a], p[b], p[c]))/2.0,
     1):
                                                                  vol += volume(P3(0, 0, 0), p[a], p[b], p[c])/6.0;
  llf sum = 0;
  int cnt = segs[0].second;
                                                               6.24 3D Projection [68f350]
  rep(j,1,sz(segs)) {
                                                               using P3F = valarray<llf>;
   if (!cnt) sum += segs[j].first - segs[j - 1].first;
                                                               P3F toP3F(P3 p) { return {p.x, p.y, p.z}; }
   cnt += segs[j].second;
                                                               llf dot(P3F a, P3F b) {
                                                               return a[0]*b[0]+a[1]*b[1]+a[2]*b[2];
  ret += cross(A,B) * sum;
                                                               P3F housev(P3 A, P3 B, int s) {
 return ret / 2;
                                                                const llf a = abs(A), b = abs(B);
                                                                return toP3F(A) / a + s * toP3F(B) / b;
6.22 3D Point [46b73b]
                                                               P project(P3 p, P3 q) {
struct P3 {
                                                                P3 o(0, 0, 1);
 lld x, y, z;
                                                                P3F u = housev(q, o, q.z > 0 ? 1 : -1);
 P3 operator^(const P3 &b) const {
                                                                auto pf = toP3F(p);
  return {y*b.z-b.y*z, z*b.x-b.z*x, x*b.y-b.x*y};
                                                                auto np = pf - 2 * u * dot(u, pf) / dot(u, u);
                                                                return P(np[0], np[1]);
 //Azimuthal angle (longitude) to x-axis. \in [-pi, pi]
                                                                // project p onto the plane q^Tx = 0
 llf phi() const { return atan2(y, x); }
                                                               6.25 3D Skew Line Nearest Point
 //Zenith angle (latitude) to the z-axis. \in [0, pi]
 llf theta() const { return atan2(sqrt(x*x+y*y),z); }
                                                               • L_1: \boldsymbol{v}_1 = \boldsymbol{p}_1 + t_1 \boldsymbol{d}_1, L_2: \boldsymbol{v}_2 = \boldsymbol{p}_2 + t_2 \boldsymbol{d}_2
                                                               \cdot n = d_1 \times d_2
P3 ver(P3 a, P3 b, P3 c) { return (b - a) ^ (c - a); }
                                                               \boldsymbol{\cdot} \ \boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}, \boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}
                                                               • c_1 = p_1 + \frac{(p_2 - p_1) \cdot n_2}{d_1 \cdot n_2} d_1, c_2 = p_2 + \frac{(p_1 - p_2) \cdot n_1}{d_2 \cdot n_1} d_2
lld volume(P3 a, P3 b, P3 c, P3 d) {
return dot(ver(a, b, c), d - a);
                                                               6.26 Delaunay [3a4ff1]
                                                               /* please ensure input points are unique */
P3 rotate_around(P3 p, llf angle, P3 axis) {
                                                               /* A triangulation such that no points will strictly
llf s = sin(angle), c = cos(angle);
                                                               inside circumcircle of any triangle. C should be big
 P3 u = normalize(axis);
                                                               enough s.t. the initial triangle contains all points */
 return u*dot(u, p)*(1-c) + p * c + cross(u, p)*s;
                                                               #define L(i) ((i)==0 ? 2 : (i)-1)
                                                               #define R(i) ((i)==2 ? 0 : (i)+1)
6.23 3D Convex Hull [01652a]
                                                               #define F3 for (int i = 0; i < 3; i++)
struct Face {
                                                               bool is_inf(P z) { return RE(z) \leftarrow -C || RE(z) \rightarrow= C; }
 int a, b, c;
                                                               bool in_cc(const array<P,3> &p, P q) {
                                                                i128 inf_det = 0, det = 0, inf_N, N;
 Face(int ta, int tb, int tc) : a(ta), b(tb), c(tc) {}
                                                                F3 {
                                                                 if (is_inf(p[i]) && is_inf(q)) continue;
auto preprocess(const vector<P3> &pt) {
 auto G = pt.begin();
                                                                 else if (is_inf(p[i])) inf_N = 1, N = -norm(q);
 auto a = find_if(all(pt), [&](P3 z) {
                                                                 else if (is_inf(q)) inf_N = -1, N = norm(p[i]);
 return z != *G; }) - G;
                                                                 else inf_N = 0, N = norm(p[i]) - norm(q);
 auto b = find_if(all(pt), [&](P3 z) {
                                                                 lld D = cross(p[R(i)] - q, p[L(i)] - q);
 return ver(*G, pt[a], z) != P3(0, 0, 0); }) - G;
auto c = find_if(all(pt), [&](P3 z) {
                                                                 inf_det += inf_N * D; det += N * D;
  return volume(*G, pt[a], pt[b], z) != 0; }) - G;
                                                                return inf_det != 0 ? inf_det > 0 : det > 0;
 vector<size_t> id;
 for (size_t i = 0; i < pt.size(); i++)</pre>
                                                               P v[maxn];
  if (i != a && i != b && i != c) id.push_back(i);
                                                               struct Tri;
 return tuple{a, b, c, id};
                                                               struct E {
                                                                Tri *t; int side;
// return the faces with pt indexes
                                                                E(Tri *t_=0, int side_=0) : t(t_), side(side_) {}
// all points coplanar case will WA
vector<Face> convex_hull_3D(const vector<P3> &pt) {
                                                               struct Tri {
 const int n = int(pt.size());
                                                                array<int,3> p; array<Tri*,3> ch; array<E,3> e;
                                                                Tri(int a=0, int b=0, int c=0) : p{a, b, c}, ch{} {}
 if (n <= 3) return {}; // be careful about edge case</pre>
                                                                bool has_chd() const { return ch[0] != nullptr; }
 vector<Face> now;
 vector<vector<int>> z(n, vector<int>(n));
                                                                bool contains(int q) const {
 auto [a, b, c, ord] = preprocess(pt);
                                                                 F3 if (ori(v[p[i]], v[p[R(i)]], v[q]) < 0)
 now.emplace_back(a, b, c); now.emplace_back(c, b, a);
                                                                  return false;
 for (auto i : ord) {
                                                                 return true:
  vector<Face> next;
  for (const auto &f : now) {
                                                                bool check(int q) const {
   lld v = volume(pt[f.a], pt[f.b], pt[f.c], pt[i]);
                                                                 return in_cc({v[p[0]], v[p[1]], v[p[2]]}, v[q]); }
   if (v <= 0) next.push_back(f);</pre>
                                                               } pool[maxn * 10], *it, *root;
   z[f.a][f.b] = z[f.b][f.c] = z[f.c][f.a] = sgn(v);
                                                               void link(const E &a, const E &b) {
                                                                if (a.t) a.t->e[a.side] = b;
                                                                if (b.t) b.t->e[b.side] = a;
  const auto F = [\&](int x, int y) \{
   if (z[x][y] > 0 && z[y][x] <= 0)
    next.emplace_back(x, y, i);
                                                               void flip(Tri *A, int a) {
                                                                auto [B, b] = A->e[a]; /* flip edge between A,B */
                                                                if (!B || !A->check(B->p[b])) return;
  for (const auto &f : now)
   F(f.a, f.b), F(f.b, f.c), F(f.c, f.a);
                                                                Tri *X = new (it++) Tri(A->p[R(a)], B->p[b], A->p[a]);
                                                                Tri *Y = new (it++) Tri(B->p[R(b)], A->p[a], B->p[b]);
  now = next;
```

```
link(E(X, 0), E(Y, 0));
 link(E(X, 1), A \rightarrow e[L(a)]); link(E(X, 2), B \rightarrow e[R(b)]);
                                                                  using P = pair<lld, int>;
 link(E(Y, 1), B\rightarrow e[L(b)]); link(E(Y, 2), A\rightarrow e[R(a)]);
                                                                  void dfs(int x, int y, P &mn, Node *r) {
A->ch = B->ch = {X, Y, nullptr};
flip(X, 1); flip(X, 2); flip(Y, 1); flip(Y, 2);
                                                                  if (!r || !touch(x, y, mn.first, r)) return;
mn = min(mn, P(dis2(r->x, r->y, x, y), r->id));
                                                                   if (r->f == 1 ? y < r->y : x < r->x)
                                                                    dfs(x, y, mn, r\rightarrow L), dfs(x, y, mn, r\rightarrow R);
void add_point(int p) {
Tri *r = root;
                                                                   else
 while (r->has_chd()) for (Tri *c: r->ch)
                                                                    dfs(x, y, mn, r\rightarrow R), dfs(x, y, mn, r\rightarrow L);
 if (c && c->contains(p)) { r = c; break; }
 array<Tri*, 3> t; /* split into 3 triangles */
                                                                  int query(int x, int y) {
 F3 t[i] = new (it++) Tri(r->p[i], r->p[R(i)], p);
                                                                   P mn(INF, -1); dfs(x, y, mn, root);
 F3 link(E(t[i], 0), E(t[R(i)], 1));
                                                                   return mn.second;
 F3 link(E(t[i], 2), r->e[L(i)]);
 r->ch = t:
                                                                } tree;
F3 flip(t[i], 2);
                                                                6.29 kd Closest Pair (3D ver.) [84d9eb]
                                                                llf solve(vector<P> v) {
auto build(const vector<P> &p) {
                                                                 shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,</pre>
 it = pool; int n = (int)p.size();
 vector<int> ord(n); iota(all(ord), 0);
                                                                  unordered_map<lld, int>>> m;
 shuffle(all(ord), mt19937(114514));
                                                                  llf d = dis(v[0], v[1]);
 root = new (it++) Tri(n, n + 1, n + 2);
                                                                  auto Idx = [\&d] (llf x) \rightarrow lld {
 copy_n(p.data(), n, v); v[n++] = P(-C, -C);
                                                                  return round(x * 2 / d) + 0.1; };
 v[n++] = P(C * 2, -C); v[n++] = P(-C, C * 2);
                                                                  auto rebuild_m = [&m, &v, &Idx](int k) {
 for (int i : ord) add_point(i);
                                                                  m.clear();
 vector<array<int, 3>> res;
                                                                   for (int i = 0; i < k; ++i)</pre>
 for (Tri *now = pool; now != it; now++)
                                                                    m[Idx(v[i].x)][Idx(v[i].y)]
  if (!now->has_chd()) res.push_back(now->p);
                                                                     [Idx(v[i].z)] = i;
 return res:
                                                                  }; rebuild_m(2);
                                                                  for (size_t i = 2; i < v.size(); ++i) {</pre>
                                                                   const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
6.27 Build Voronoi [94f000]
                                                                      kz = Idx(v[i].z); bool found = false;
void build_voronoi_cells(auto &&p, auto &&res) {
                                                                   for (int dx = -2; dx \le 2; ++dx) {
 vector<vector<int>> adj(p.size());
                                                                    const lld nx = dx + kx;
 for (auto f: res) F3 {
                                                                    if (m.find(nx) == m.end()) continue;
  int a = f[i], b = f[R(i)];
if (a >= p.size() || b >= p.size()) continue;
                                                                    auto& mm = m[nx];
  adj[a].emplace_back(b);
                                                                    for (int dy = -2; dy <= 2; ++dy) {
                                                                     const lld ny = dy + ky;
// use `adj` and `p` and HPI to build cells
for (size_t i = 0; i < p.size(); i++) {</pre>
                                                                     if (mm.find(ny) == mm.end()) continue;
                                                                     auto& mmm = mm[ny];
                                                                     for (int dz = -2; dz \le 2; ++dz) {
  vector<Line> ls = frame; // the frame
                                                                      const lld nz = dz + kz;
  for (int j : adj[i]) {
                                                                      if (mmm.find(nz) == mmm.end()) continue;
  P m = p[i] + p[j], d = rot90(p[j] - p[i]);
                                                                      const int p = mmm[nz];
   assert (norm(d) != 0);
                                                                      if (dis(v[p], v[i]) < d) {</pre>
   ls.emplace_back(m, m + d); // doubled coordinate
                                                                       d = dis(v[p], v[i]);
  } // HPI(ls)
                                                                       found = true;
}
        kd Tree (Nearest Point) [f733e5]
                                                                    }
struct KDTree {
 struct Node {
                                                                   if (found) rebuild_m(i + 1);
  int x, y, x1, y1, x2, y2, id, f; Node *L, *R;
                                                                   else m[kx][ky][kz] = i;
 } tree[maxn], *root;
 lld dis2(int x1, int y1, int x2, int y2) {
lld dx = x1 - x2, dy = y1 - y2;
                                                                  return d;
  return dx * dx + dy * dy;
                                                                6.30 Simulated Annealing [4e0fe5]
                                                                llf anneal() {
 static bool cmpx(Node& a, Node& b) { return a.x<b.x; }</pre>
                                                                  mt19937 rnd_engine(seed);
 static bool cmpy(Node& a, Node& b) { return a.y<b.y; }</pre>
                                                                  uniform_real_distribution<llf> rnd(0, 1);
 void init(vector<pair<int,int>> &ip) {
                                                                  const llf dT = 0.001;
  for (int i = 0; i < ssize(ip); i++)</pre>
                                                                    'Argument p
   tie(tree[i].x, tree[i].y) = ip[i], tree[i].id = i;
                                                                  Ilf S_cur = calc(p), S_best = S_cur;
for (llf T = 2000; T > EPS; T -= dT) {
  root = build(0, (int)ip.size()-1, 0);
                                                                  // Modify p to p_prime
const llf S_prime = calc(p_prime);
 Node* build(int L, int R, int d) {
  if (L>R) return nullptr;
                                                                   const llf delta_c = S_prime - S_cur;
  int M = (L+R)/2;
                                                                   llf prob = min((llf)1, exp(-delta_c / T));
  nth_element(tree+L,tree+M,tree+R+1,d%2?cmpy:cmpx);
                                                                   if (rnd(rnd_engine) <= prob)</pre>
  Node &o = tree[M]; o.f = d \% 2;
                                                                    S_cur = S_prime, p = p_prime;
 o.x1 = o.x2 = o.x; o.y1 = o.y2 = o.y;
                                                                   if (S_prime < S_best) // find min</pre>
  o.L = build(L, M-1, d+1); o.R = build(M+1, R, d+1);
                                                                    S_best = S_prime, p_best = p_prime;
  for (Node *s: {o.L, o.R}) if (s) {
   o.x1 = min(o.x1, s->x1); o.x2 = max(o.x2, s->x2);
                                                                  return S_best;
   o.y1 = min(o.y1, s->y1); o.y2 = max(o.y2, s->y2);
                                                                6.31 Triangle Centers [adb146]
  return tree+M;
                                                                0 = ... // see min circle cover
 bool touch(int x, int y, lld d2, Node *r){
                                                                G = (A + B + C) / 3;
                                                                H = G * 3 - 0 * 2; // orthogonal center
llf a = abs(B - C), b = abs(A - C), c = abs(A - B);
  lld d = (lld) sqrt(d2) +1;
  return x >= r->x1 - d && x <= r->x2 + d &&
         y >= r->y1 - d && y <= r->y2 + d;
                                                                I = (a * A + b * B + c * C) / (a + b + c);
```

```
// FermatPoint: minimizes sum of distance
                                                                 struct exSAM {
// if max. angle >= 120 deg then vertex
                                                                  int len[maxn * 2], link[maxn * 2]; // maxlen, suflink
// otherwise, make eq. triangle AB'C, CA'B, BC'A
// line AA', BB', CC' intersects at P
7 Stringology
                                                                  int next[maxn * 2][maxc], tot; // [0, tot), root = 0
                                                                  int ord[maxn * 2]; // topo. order (sort by length)
int cnt[maxn * 2]; // occurence
                                                                  int newnode() {
      Hash [ce7fad]
                                                                   fill_n(next[tot], maxc, 0);
template <int P = 127, int Q = 1051762951>
                                                                   return len[tot] = cnt[tot] = link[tot] = 0, tot++;
class Hash {
 vector<int> h, p;
                                                                  void init() { tot = 0, newnode(), link[0] = -1; }
int insertSAM(int last, int c) {
public:
 Hash(const auto &s) : h(s.size()+1), p(s.size()+1) {
                                                                   int cur = next[last][c];
  for (size_t i = 0; i < s.size(); ++i)</pre>
                                                                   len[cur] = len[last] + 1;
  h[i + 1] = add(mul(h[i], P), s[i]);
                                                                   int p = link[last];
  generate(all(p), [x = 1, y = 1, this]() mutable {
                                                                   while (p != -1 && !next[p][c])
   return y = x, x = mul(x, P), y; });
                                                                    next[p][c] = cur, p = link[p];
                                                                   if (p == -1) return link[cur] = 0, cur;
 int query(int l, int r) const { // 0-base [l, r)
                                                                   int q = next[p][c];
  return sub(h[r], mul(h[l], p[r - l]));
                                                                   if (len[p] + 1 == len[q]) return link[cur] = q, cur;
                                                                   int clone = newnode();
};
7.2
                                                                   for (int i = 0; i < maxc; ++i)</pre>
      Suffix Array [e9e77d]
                                                                    next[clone][i] = len[next[q][i]] ? next[q][i] : 0;
auto sais(const auto &s) {
                                                                   len[clone] = len[p] + 1;
 const int n = (int)s.size(), z = ranges::max(s) + 1;
                                                                   while (p != -1 && next[p][c] == q)
 vector<int> c(z); for (int x : s) ++c[x];
                                                                    next[p][c] = clone, p = link[p];
 partial_sum(all(c), begin(c));
                                                                   link[link[cur] = clone] = link[q];
 vector<int> sa(n); auto I = ranges::iota_view(0, n);
                                                                   link[q] = clone;
 if (ranges::max(c) <= 1) {</pre>
                                                                   return cur;
  for (int i : I) sa[--c[s[i]]] = i;
  return sa;
                                                                  void insert(const string &s) {
                                                                   int cur = 0;
 vector<bool> t(n); t[n - 1] = true;
                                                                   for (char ch : s) {
 for (int i = n - 2; i >= 0; --i)
t[i] = (s[i]==s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);</pre>
                                                                    int &nxt = next[cur][int(ch - 'a')];
                                                                    if (!nxt) nxt = newnode();
 auto is_lms = ranges::views::filter([&t](int x) {
                                                                    cnt[cur = nxt] += 1;
  return x && t[x] && !t[x - 1]; });
                                                                   }
 const auto induce = [&] {
  for (auto x = c; int y : sa)
                                                                  void build() {
   if (y--) if (!t[y]) sa[x[s[y] - 1]++] = y;
                                                                   queue<int> q; q.push(0);
  for (auto x = c; int y : sa | views::reverse)
                                                                   while (!q.empty()) {
   if (y--) if (t[y]) sa[--x[s[y]]] = y;
                                                                    int cur = q.front(); q.pop();
 };
                                                                    for (int i = 0; i < maxc; ++i)</pre>
 vector<int> lms, q(n); lms.reserve(n);
                                                                     if (next[cur][i]) q.push(insertSAM(cur, i));
 for (auto x = c; int i : I | is_lms) {
  q[i] = int(lms.size());
                                                                   vector<int> lc(tot);
for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
  lms.push_back(sa[--x[s[i]]] = i);
                                                                   partial_sum(all(lc), lc.begin());
 induce(); vector<int> ns(lms.size());
                                                                   for (int i = 1; i < tot; ++i) ord[--lc[len[i]]] = i;</pre>
 for (int j = -1, nz = 0; int i : sa | is_lms) {
  if (j >= 0) {
                                                                  void solve() {
   int len = min({n - i, n - j, lms[q[i] + 1] - i});
                                                                   for (int i = tot - 2; i >= 0; --i)
   ns[q[i]] = nz += lexicographical_compare(
                                                                    cnt[link[ord[i]]] += cnt[ord[i]];
     begin(s) + j, begin(s) + j + len,
     begin(s) + i, begin(s) + i + len);
                                                                };
                                                                 7.4 KMP [281185]
  j = i;
                                                                 vector<int> kmp(const auto &s) {
 ranges::fill(sa, 0); auto nsa = sais(ns);
for (auto x = c; int y : nsa | views::reverse)
                                                                  vector<int> f(s.size());
                                                                  for (int i = 1, k = 0; i < (int)s.size(); ++i) {
  while (k > 0 && s[i] != s[k]) k = f[k - 1];
  y = lms[y], sa[--x[s[y]]] = y;
 return induce(), sa;
                                                                   f[i] = (k += (s[i] == s[k]));
// sa[i]: sa[i]-th suffix is the
                                                                  return f;
// i-th lexicographically smallest suffix.
// hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
                                                                 vector<int> search(const auto &s, const auto &t) {
                                                                  // return 0-indexed occurrence of t in s
struct Suffix {
 int n; vector<int> sa, hi, rev;
                                                                  vector<int> f = kmp(t), r;
                                                                  for (int i = 0, k = 0; i < (int)s.size(); ++i) {
  while (k > 0 && s[i] != t[k]) k = f[k - 1];
 Suffix(const auto &s) : n(int(s.size())),
 hi(n), rev(n) {
  vector<int> _s(n + 1); _s[n] = 0;
                                                                   k += (s[i] == t[k]);
  copy(all(s), begin(_s)); // s shouldn't contain 0
                                                                   if (k == (int)t.size()) {
  sa = sais(_s); sa.erase(sa.begin());
                                                                    r.push_back(i - t.size() + 1);
  for (int i = 0; i < n; ++i) rev[sa[i]] = i;</pre>
                                                                    k = f[k - 1];
  for (int i = 0, h = 0; i < n; ++i) {
  if (!rev[i]) { h = 0; continue; }</pre>
                                                                   }
                                                                  }
   for (int j = sa[rev[i] - 1]; i + h < n && j + h < n</pre>
                                                                  return r;
     && s[i + h] == s[j + h];) ++h;
   hi[rev[i]] = h ? h-- : 0;
                                                                 7.5 Z value [6a7fd0]
  }
                                                                vector<int> Zalgo(const string &s) {
 }
                                                                  vector<int> z(s.size(), s.size());
for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {</pre>
     Ex SAM [58374b]
                                                                   int j = clamp(r - i, 0, z[i - l]);
```

```
for (; i + j < z[0] \text{ and } s[i + j] == s[j]; ++j);
  if (i + (z[i] = j) > r) r = i + z[l = i];
}
7.6 Manacher [c938a9]
vector<int> manacher(const string &S) {
const int n = (int)S.size(), m = n * 2 + 1;
vector<int> z(m);
string t = "."; for (char c: S) t += c, t += '.';
for (int i = 1, l = 0, r = 0; i < m; ++i) {
z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
 while (i - z[i] >= 0 \&\& i + z[i] < m) {
   if (t[i - z[i]] == t[i + z[i]]) ++z[i];
   else break:
 if (i + z[i] > r) r = i + z[i], l = i;
return z; // the palindrome lengths are z[i] - 1
/* for (int i = 1; i + 1 < m; ++i) {
 int l = (i - z[i] + 2) / 2, r = (i + z[i]) / 2;
 if (l != r) // [l, r) is maximal palindrome
     Lyndon Factorization [d22cc9]
// partition s = w[0] + w[1] + ... + w[k-1],
// w[0] >= w[1] >= ... >= w[k-1]
// each w[i] strictly smaller than all its suffix
void duval(const auto &s, auto &&report) {
for (int n = (int)s.size(), i = 0, j, k; i < n; ) {</pre>
  for (j = i + 1, k = i; j < n && s[k] <= s[j]; j++)
  k = (s[k] < s[j] ? i : k + 1);
  // if (i < n / 2 && j >= n / 2) {
 // for min cyclic shift, call duval(s + s)
 // then here s.substr(i, n / 2) is min cyclic shift
  for (; i <= k; i += j - k)
  report(i, j - k); // s.substr(l, len)
} // tested @ luogu 6114, 1368 & UVA 719
7.8 Main Lorentz [615b8f]
vector<pair<int, int>> rep[kN]; // 0-base [l, r]
void main_lorentz(const string &s, int sft = 0) {
const int n = s.size();
if (n == 1) return;
const int nu = n / 2, nv = n - nu;
const string u = s.substr(0, nu), v = s.substr(nu);
    ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend());
main_lorentz(u, sft), main_lorentz(v, sft + nu);
auto get_z = [](const vector<int> &z, int i) {
  return (0 <= i and i < (int)z.size()) ? z[i] : 0; };</pre>
auto add_rep = [&](bool left, int c, int l, int k1,
    int k2) {
  const int L = max(1, l - k2), R = min(l - left, k1);
  if (L > R) return;
  if (left) rep[l].emplace_back(sft + c - R, sft + c -
    L);
  else rep[l].emplace_back(sft + c - R - l + 1, sft + c
     - L - l + 1);
 for (int cntr = 0; cntr < n; cntr++) {</pre>
  int l, k1, k2;
 if (cntr < nu) {
  l = nu - cntr;</pre>
  k1 = get_z(z1, nu - cntr);
  k2 = get_z(z2, nv + 1 + cntr);
  } else {
   l = cntr - nu + 1;
  k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
   k2 = get_z(z4, (cntr - nu) + 1);
 if (k1 + k2 >= 1)
   add_rep(cntr < nu, cntr, l, k1, k2);</pre>
}
7.9
    BWT [5a9b3a]
vector<int> v[SIGMA];
void BWT(char *ori, char *res) {
// make ori -> ori + ori
```

```
void iBWT(char *ori, char *res) {
 for (int i = 0; i < SIGMA; i++) v[i].clear();</pre>
 const int len = strlen(ori);
 for (int i = 0; i < len; i++)</pre>
  v[ori[i] - 'a'].push_back(i);
 vector<int> a;
 for (int i = 0, ptr = 0; i < SIGMA; i++)</pre>
  for (int j : v[i]) {
   a.push_back(j);
   ori[ptr++] = 'a' + i;
 for (int i = 0, ptr = 0; i < len; i++) {</pre>
  res[i] = ori[a[ptr]];
  ptr = a[ptr];
 res[len] = 0;
7.10 Palindromic Tree [0673ee]
struct PalindromicTree {
 struct node {
  int nxt[26], f, len; // num = depth of fail link
                   // = #pal_suffix of this node
  int cnt, num;
  node(int l = 0) : nxt{}, f(0), len(l), cnt(0), num(0)
 };
 vector<node> st; vector<char> s; int last, n;
 void init() {
  st.clear(); s.clear();
  last = 1; n = 0;
  st.push_back(0); st.push_back(-1);
  st[0].f = 1; s.push_back(-1);
 int getFail(int x) {
  while (s[n - st[x].len - 1] != s[n]) x = st[x].f;
  return x;
 void add(int c) {
  s.push_back(c -= 'a'); ++n;
  int cur = getFail(last);
  if (!st[cur].nxt[c]) {
   int now = st.size();
   st.push_back(st[cur].len + 2);
   st[now].f = st[getFail(st[cur].f)].nxt[c];
   st[cur].nxt[c] = now;
   st[now].num = st[st[now].f].num + 1;
  last = st[cur].nxt[c]; ++st[last].cnt;
 void dpcnt() { // cnt = #occurence in whole str
  for (int i = st.size() - 1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
 int size() { return st.size() - 2; }
} pt;
/* usaae
string s; cin >> s; pt.init();
for (int i = 0; i < SZ(s); i++) {
 int prvsz = pt.size(); pt.add(s[i]);
 if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
  // pal @ [l,r]: s.substr(l, r-l+1)
 }
} */
8
     Misc
8.1
     Theorems
Spherical Coordinate
                                           r = \sqrt{x^2 + y^2 + z^2}
 x=r\sin\theta\cos\phi
 y = r \sin \theta \sin \phi
                                    \theta = \mathrm{acos}(z/\sqrt{x^2 + y^2 + z^2})
```

// then build suffix array

Spherical Cap

 $z=r\cos\theta$

- · A portion of a sphere cut off by a plane.
- r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ :

 $\phi = \mathsf{atan2}(u, x)$

- $\begin{array}{l} \cdot \text{ Volume} = \pi h^2 (3r-h)/3 = \pi h (3a^2+h^2)/6 = \pi r^3 (2+\cos\theta)(1-\cos\theta)^2/3. \\ \cdot \text{ Area} = 2\pi r h = \pi (a^2+h^2) = 2\pi r^2 (1-\cos\theta). \end{array}$

Sherman-Morrison formula

 $(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1+v^{\mathsf{T}}A^{-1}u}$

Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} =$ d(i), $L_{ij} = -c$ where c is the number of edge (i, j) in \tilde{G} .

- The number of undirected spanning in G is $\det(\tilde{L}_{11})$.
- The number of directed spanning tree rooted at r in G is $\det(\tilde{L}_{rr})$.

BEST Theorem

#{Eulerian circuits} = #{arborescences rooted at 1} $\cdot \prod_{v \in V} (\deg(v) - 1)!$

Random Walk on Graph

Let P be the transition matrix of a strongly connected directed graph, $\sum_{i} P_{i,j} = 1$. Let $F_{i,j}$ be the expected time to reach j from i. Let g_i be the expected time from i to i, G = diag(g) and J be a matrix all of 1, i.e. $J_{i,j} = 1$.

First solve G: let $\pi P = \pi$ be a stationary distribution. Then $\pi_i g_i = 1$. The rank of I-P is n-1, so we first solve a special solution X such that (I-P)X=J-G and adjust X to F by $F_{i,j}=X_{i,j}-X_{j,j}$.

Tutte Matrix

For i < j, $d_{ij} = x_{ij}$ (in practice, a random number) if $(i,j) \in E$, otherwise $d_{ij}=0$. For $i \geq j, d_{ij}=-d_{ji}$. $\frac{\operatorname{rank}(D)}{2}$ is the maximum matching.

Cayley's Formula

- · Given a degree sequence d_1, d_2, \ldots, d_n for each labeled vertices, there're $rac{(n-2)!}{-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1, 2, \ldots, k$ belong to different components. Then $T_{n,k} =$

Erdős-Gallai theorem

A sequence of non-negative integers $d_1 > d_2 > \ldots > d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and $\sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k)$ holds for all 1 < k < n.

Havel-Hakimi algorithm

Find the vertex who has greatest degree unused, connect it with other greatest vertex.

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \ge \cdots \ge a_n$ and b_1, \ldots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \le \sum_{i=1}^n \min(b_i, k)$ holds for every $1 \le k \le n$.

Fulkerson–Chen–Anstee theorem

A sequence $(a_1,b_1),\ldots,(a_n,b_n)$ of nonnegative integer pairs with $a_1\geq \cdots \geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq a_i$ $\sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$ **Euler's planar graph formula** $\underbrace{V-E+F}_{} = C+1. \ E \leq 3V-6 \text{ (when } V \geq 3\text{)}$

Pick's theorem

For simple polygon, when points are all integer, we have \boldsymbol{A} #{lattice points in the interior} $+\frac{1}{2}$ #{lattice points on the boundary} -1

Matroid Intersection

Given matroids $M_1=(G,I_1), M_2=(G,I_2)$, find maximum $S\in I_1\cap I_2$. For each iteration, build the directed graph and find a shortest path from s to t.

```
• s \rightarrow x : S \sqcup \{x\} \in I_1
```

• $x \to t: S \sqcup \{x\} \in I_2$ • $y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1$ (y is in the unique circuit of $S \sqcup \{x\}$) • $x \to y: S \setminus \{y\} \sqcup \{x\} \in I_2$ (y is in the unique circuit of $S \sqcup \{x\}$)

Alternate the path, and |S| will increase by 1. In each iteration, |E| = O(RN), where $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)), N=|G|$. For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Find the shortest path by Bellman-Ford. The maximum iteration of Bellman-Ford is 2R+1.

Dual of LP

Primal	Dual
Maximize $c^{T}x$ s.t. $Ax \leq b$, $x \geq 0$	Minimize $b^{T}y$ s.t. $A^{T}y \geq c$, $y \geq 0$
Maximize $c^{T}x$ s.t. $Ax \leq b$	Minimize $b^{T}y$ s.t. $A^{T}y = c$, $y \geq 0$
Maximize $c^{T}x$ s.t. $Ax = b$, $x \geq 0$	Minimize $b^{T}y$ s.t. $A^{T}y \geq c$

Minimax Theorem

Let $f: X \times Y \to \mathbb{R}$ be continuous where $X \subseteq \mathbb{R}^n, Y \subseteq \mathbb{R}^m$ are compact and convex. If $f(\cdot,y):X\to\mathbb{R}$ is concave for fixed y, and $f(x,\cdot):Y\to\mathbb{R}$ is convex for fixed x, then $\max_{x \in X} \min_{y \in Y} f(x, y) = \min_{y \in Y} \max_{x \in X} f(x, y)$, e.g. $f(x, y) = x^{\mathsf{T}} A y$ for

zero-sum matrix game. Parallel Axis Theorem

The second moment of area is $I_z=\iint x^2+y^2\mathrm{d}A.$ $I_{z'}=I_z+Ad^2$ where d is the distance between two parallel axis z,z'.

Weight Matroid Intersection [d00ee8] 8.2

```
struct Matroid {
Matroid(bitset<N>); // init from an independent set
bool can_add(int); // check if break independence
Matroid remove(int); // removing from the set
auto matroid_intersection(const vector<int> &w) {
const int n = (int)w.size(); bitset<N> S;
for (int sz = 1; sz <= n; sz++) {</pre>
 Matroid M1(S), M2(S); vector<vector<pii>>> e(n + 2);
 for (int j = 0; j < n; j++) if (!S[j]) {</pre>
  if (M1.can_add(j)) e[n].eb(j, -w[j]);
  if (M2.can_add(j)) e[j].eb(n + 1, 0);
  for (int i = 0; i < n; i++) if (S[i]) {</pre>
  Matroid T1 = M1.remove(i), T2 = M2.remove(i);
```

```
for (int j = 0; j < n; j++) if (!S[j]) {</pre>
    if (T1.can_add(j)) e[i].eb(j, -w[j]);
    if (T2.can_add(j)) e[j].eb(i, w[i]);
  } // maybe implicit build graph for more speed
  vector<pii> d(n + 2, \{INF, 0\}); d[n] = \{0, 0\};
  vector<int> prv(n + 2, -1);
  // change to SPFA for more speed, if necessary
  for (int upd = 1; upd--; )
for (int u = 0; u < n + 2; u++)</pre>
     for (auto [v, c] : e[u]) {
      pii x(d[u].first + c, d[u].second + 1);
      if (x < d[v]) d[v] = x, prv[v] = u, upd = 1;
  if (d[n + 1].first >= INF) break;
  for (int x = prv[n+1]; x!=n; x = prv[x]) S.flip(x);
  // S is the max-weighted independent set w/ size sz
 return S;
} // from Nacl
8.3 Stable Marriage
1: Initialize m \in M and w \in W to free
  while \exists free man m who has a woman w to propose to do
        \leftarrow first woman on m's list to whom m has not yet proposed
     if \exists some pair (m', w) then
        if w prefers m to m^\prime then
          m' \gets \mathit{free}
6:
           (\underline{m},w) \gets \mathsf{engaged}
8:
        end if
9:
10:
        (m,w) \leftarrow \mathsf{engaged}
11.
     end if
12: end while
8.4 Bitset LCS [4155ab]
cin >> n >> m;
for (int i = 1, x; i <= n; ++i)</pre>
```

```
cin >> x, p[x].set(i);
for (int i = 1, x; i <= m; ++i) {</pre>
 cin >> x, (g = f) = p[x];
 f.shiftLeftByOne(), f.set(0);
 ((f = g - f) ^= g) \&= g;
cout << f.count() << '\n';</pre>
```

Prefix Substring LCS [7d8faf] 8.5

```
void all_lcs(string S, string T) { // 0-base
 vector<size_t> h(T.size()); iota(all(h), 1);
 for (size_t a = 0; a < S.size(); ++a) {</pre>
  for (size_t c = 0, v = 0; c < T.size(); ++c)</pre>
   if (S[a] == T[c] || h[c] < v) swap(h[c], v);</pre>
  // here, LCS(s[0, a], t[b, c]) =
  // c - b + 1 - sum([h[i] > b] | i <= c)
} // test @ yosupo judge
```

8.6 Convex 1D/1D DP [e5ab4b]

return dp;

```
struct S { int i, l, r; };
auto solve(int n, int k, auto &w) {
 vector<int64_t> dp(n + 1);
 auto f = [\&](int l, int r) \rightarrow int64_t {
  if (r - l > k) return -INF;
  return dp[l] + w(l + 1, r);
 dp[0] = 0;
 deque<S> dq; dq.emplace_back(0, 1, n);
 for (int i = 1; i <= n; ++i) {</pre>
  dp[i] = f(dq.front().i, i);
  while (!dq.empty() && dq.front().r <= i)</pre>
   dq.pop_front();
  dq.front().l = i + 1;
  while (!dq.empty() &&
     f(i, dq.back().l) >= f(dq.back().i, dq.back().l))
   dq.pop_back();
  int p = i + 1;
  if (!dq.empty()) {
   auto [j, l, r] = dq.back();
   for (int s = 1 << 20; s; s >>= 1)
if (l + s <= n && f(i, l + s) < f(j, l + s))</pre>
      l += s:
   dq.back().r = l; p = l + 1;
  if (p <= n) dq.emplace_back(i, p, n);</pre>
```

```
} // test @ tioj 烏龜疊疊樂
                                                              ret.push_back(2);
                                                              for (int i = 5; i <= n; i += 2) ret.push_back(i);</pre>
8.7 ConvexHull Optimization [b4318e]
                                                              ret.push_back(1); ret.push_back(3);
struct L {
                                                             } else {
 mutable lld a, b, p;
                                                              for (int i = 2; i <= n; i += 2) ret.push_back(i);</pre>
 bool operator<(const L &r) const {</pre>
                                                              for (int i = 1; i <= n; i += 2) ret.push_back(i);</pre>
  return a < r.a; /* here */ }
 bool operator<(lld x) const { return p < x; }</pre>
lld Div(lld a, lld b) {
  return a / b - ((a ^ b) < 0 && a % b); }</pre>
                                                            8.12 Tree Knapsack [f42766]
                                                            vector<int> G[N]; int dp[N][K]; pair<int,int> obj[N];
                                                            void dfs(int u, int mx) {
struct DynamicHull : multiset<L, less<>>> {
 static const lld kInf = 1e18;
                                                             for (int s : G[u]) {
 bool Isect(iterator x, iterator y) {
                                                              auto [w, v] = obj[s];
                                                              if (mx < w) continue;</pre>
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a)
                                                              for (int i = 0; i <= mx - w; i++)</pre>
   x->p = x->b > y->b ? kInf : -kInf; /* here */
                                                              dp[s][i] = dp[u][i];
  else x->p = Div(y->b - x->b, x->a - y->a);
                                                              dfs(s, mx - w);
                                                              for (int i = w; i <= mx; i++)</pre>
  return x->p >= y->p;
                                                               dp[u][i] = max(dp[u][i], dp[s][i - w] + v);
 void Insert(lld a, lld b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
                                                            8.13
                                                                   Manhattan MST [1008bc]
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
                                                            vector<array<int, 3>> manhattanMST(vector<P> ps) {
  while ((y = x) != begin() && (--x)->p >= y->p)
                                                             vector<int> id(ps.size()); iota(all(id), 0);
   Isect(x, erase(y));
                                                             vector<array<int, 3>> edges;
                                                             for (int k = 0; k < 4; k++) {
 lld Query(lld x) { // default chmax
                                                              sort(all(id), [&](int i, int j) {
  auto l = *lower_bound(x); // to chmin:
                                                              return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y; });</pre>
  return l.a * x + l.b;
                         // modify the 2 "<>"
                                                              map<int, int> sweep;
                                                              for (int i : id) {
                                                               for (auto it = sweep.lower_bound(-ps[i].y);
8.8
      Min Plus Convolution [464dcd]
                                                                  it != sweep.end(); sweep.erase(it++))
                                                                if (P d = ps[i] - ps[it->second]; d.y > d.x) break;
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector<int> min_plus_convolution(auto &a, auto &b) {
                                                                else edges.push_back({d.y + d.x, i, it->second});
 const int n = (int)a.size(), m = (int)b.size();
 vector<int> c(n + m - 1, numeric_limits<int>::max());
                                                               sweep[-ps[i].y] = i;
 auto dc = [&](auto Y, int l, int r, int jl, int jr) {
                                                              for (P &p : ps)
  if (l > r) return;
                                                               if (k \& 1) p.x = -p.x;
  int mid = (l + r) / 2, from = -1, &best = c[mid];
  for (int j = jl; j <= jr; j++)
if (int i = mid - j; i >= 0 && i < n)</pre>
                                                               else swap(p.x, p.y);
    if (best > a[i]+b[j]) best = a[i]+b[j], from = j;
                                                             return edges; // [{w, i, j}, ...]
  Y(Y, l, mid-1, jl, from); Y(Y, mid+1, r, from, jr);
                                                            } // test @ yosupo judge
};
                                                            8.14 Binary Search On Fraction [765c5a]
 return dc(dc, 0, n-1+m-1, 0, m-1), c;
                                                            struct Q {
                                                             ll p, q;
8.9
     De-Bruijn [aa7700]
                                                             Q go(Q b, ll d) { return {p + b.p*d, q + b.q*d}; }
vector<int> de_bruijn(int k, int n) {
 // return cyclic string of len k^n s.t. every string
                                                            bool pred(Q);
 // of len n using k char appears as a substring.
                                                            // returns smallest p/q in [lo, hi] such that
 vector<int> aux(n + 1), res;
                                                            // pred(p/q) is true, and 0 <= p,q <= N
 auto db = [&](auto self, int t, int p) -> void {
                                                            Q frac_bs(ll N) {
  if (t <= n)
                                                             Q lo{0, 1}, hi{1, 0};
   for (int i = aux[t - p]; i < k; ++i, p = t)</pre>
                                                             if (pred(lo)) return lo;
    aux[t] = i, self(self, t + 1, p);
                                                             assert(pred(hi));
  else if (n % p == 0) for (int i = 1; i <= p; ++i)
                                                             bool dir = 1, L = 1, H = 1;
   res.push_back(aux[i]);
                                                             for (; L || H; dir = !dir) {
                                                              ll len = 0, step = 1;
};
                                                              for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
 return db(db, 1, 1), res;
                                                               if (Q mid = hi.go(lo, len + step);
                                                                 mid.p > N || mid.q > N || dir ^ pred(mid))
8.10 Josephus Problem [7f9ceb]
lld f(lld n, ild m, lld k) { // n people kill m for
                                                               else len += step;
    each turn
                                                              swap(lo, hi = hi.go(lo, len));
 lld s = (m - 1) \% (n - k); // O(k)
                                                              (dir ? L : H) = !!len;
 for (lld i = n - k + 1; i <= n; i++) s = (s + m) % i;
                                                             return dir ? hi : lo;
lld kth(lld n, lld m, i128 k) { // died at kth
                                                                  Barrett Reduction [d44617]
                        // O(m log(n))
 if (m == 1) return k;
 for (k = k*m+m-1; k \ge n; k = k-n + (k-n)/(m-1));
                                                            struct FastMod {
 return k;
                                                             using Big = __uint128_t; llu b, m;
                                                             FastMod(llu b) : b(b), m(-1ULL / b) {}
llu reduce(llu a) { // a % b
} // k and result are 0-based, test @ CF 101955
      N Queens Problem [31f83e]
8.11
                                                              llu r = a - (llu)((Big(m) * a) >> 64) * b;
void solve(VI &ret, int n) { // no sol when n=2,3
                                                              return r >= b ? r - b : r;
 if (n % 6 == 2) {
 for (int i = 2; i <= n; i += 2) ret.push_back(i);</pre>
  ret.push_back(3); ret.push_back(1);
 for (int i = 7; i <= n; i += 2) ret.push_back(i);</pre>
                                                            8.16 Montgomery [648fb3]
                                                            struct Mont { // Montgomery multiplication
  ret.push_back(5);
 } else if (n % 6 == 3) {
                                                             constexpr static int W = 64, L = 6;
  for (int i = 4; i <= n; i += 2) ret.push_back(i);</pre>
                                                            llu mod, R1, R2, xinv;
```

```
void set_mod(llu _mod) {
   mod = _mod; assert(mod & 1); xinv = 1;
   for (int j = 0; j < L; j++) xinv *= 2 - xinv * mod;
   assert(xinv * mod == 1);
   const u128 R = (u128(1) << W) % mod;
   R1 = llu(R); R2 = llu(R*R % mod);
}
llu redc(llu a, llu b) const {
   u128 T = u128(a) * b, m = -llu(T) * xinv;
   T += m * mod; T >>= W;
   return llu(T >= mod ? T - mod : T);
}
llu from(llu x) const {
   assert(x < mod); return redc(x, R2); }
llu get(llu a) const { return redc(a, 1); }
llu one() const { return R1; }
} mont;
// a * b % mod == get(redc(from(a), from(b)))</pre>
```