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1
Basic
```

1.1 vimrc

```
se is nu rnu bs=2 ru mouse=a encoding=utf-8
se cin et ts=4 sw=4 sts=4 t_Co=256
syn on
colorscheme ron
filetype indent on
map <F8> <ESC>:w<CR>:!clear && g++ "%" -o "%<" -
    fsanitize=address -fsanitize=undefined -g && echo
    success<CR>
map <F9> <ESC>:w<CR>:!clear && g++ "%" -o "%<" -02 &&
    echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
```

1.2 Increase Stack

```
const int size = 256 << 20;</pre>
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));
```

1.3 Pragma Optimization

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,sse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
```

1.4 IO Optimization

```
static inline int gc() {
 static char buf[ 1 << 20 ], *p = buf, *end = buf;
 if ( p == end ) {
  end = buf + fread( buf, 1, 1 << 20, stdin );
   if ( end == buf ) return EOF;
  p = buf;
 return *p++;
template < typename T >
static inline bool gn( T &_ ) {
 register int c = gc(); register T __ = 1; _ = 0;
while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
if(c == '-') { __ = -1; c = gc(); }
if(c == EOF) return false;
 while('0'<=c&&c<='9') _{-} = _{-} * 10 + c _{-} '0', c = gc();
 _ *= __;
 return true;
template < typename T, typename ...Args >
static inline bool gn( T &x, Args &...args )
{ return gn(x) && gn(args...); }
```

2 Data Structure

```
2.1 Bigint
class BigInt{
private
using lld = int_fast64_t;
#define PRINTF_ARG PRIdFAST64
#define LOG_BASE_STR "9"
static constexpr lld BASE = 1000000000;
static constexpr int LOG_BASE = 9;
vector<lld> dig; bool neg;
inline int len() const { return (int) dig.size(); }
inline int cmp_minus(const BigInt& a) const {
 if(len() == 0 && a.len() == 0) return 0;
 if(neg ^ a.neg)return a.neg ^ 1;
 if(len()!=a.len())
   return neg?a.len()-len():len()-a.len();
 for(int i=len()-1;i>=0;i--) if(dig[i]!=a.dig[i])
  return neg?a.dig[i]-dig[i]:dig[i]-a.dig[i];
 return 0;
inline void trim(){
 while(!dig.empty()&&!dig.back())dig.pop_back();
 if(dig.empty()) neg = false;
public:
BigInt(): dig(vector<lld>()), neg(false){}
BigInt(lld a): dig(vector<lld>()){
 neg = a<0; dig.push_back(abs(a));</pre>
 trim();
BigInt(const string& a): dig(vector<lld>()){
 assert(!a.empty()); neg = (a[0]=='-');
 for(int i=((int)a.size())-1;i>=neg;i-=LOG_BASE){
  11d cur = 0;
   for(int j=min(LOG_BASE-1,i-neg);j>=0;j--)
   cur = cur*10+a[i-j]-'0';
  dig.push_back(cur);
 } trim();
inline bool operator<(const BigInt& a)const
 {return cmp_minus(a)<0;}
inline bool operator<=(const BigInt& a)const</pre>
 {return cmp_minus(a)<=0;}
inline bool operator==(const BigInt& a)const
  {return cmp_minus(a)==0;}
 inline bool operator!=(const BigInt& a)const
  {return cmp_minus(a)!=0;}
inline bool operator>(const BigInt& a)const
 {return cmp_minus(a)>0;}
inline bool operator>=(const BigInt& a)const
  {return cmp_minus(a)>=0;}
BigInt operator-() const {
 BigInt ret = *this;
 ret.neg ^= 1; return ret;
BigInt operator+(const BigInt& a) const {
 if(neg) return -(-(*this)+(-a));
  if(a.neg) return (*this)-(-a);
  int n = max(a.len(), len());
 BigInt ret; ret.dig.resize(n);
 11d pro = 0;
 for(int i=0;i<n;i++) {</pre>
  ret.dig[i] = pro;
  if(i < a.len()) ret.dig[i] += a.dig[i];</pre>
  if(i < len()) ret.dig[i] += dig[i];</pre>
  pro = 0
   if(ret.dig[i] >= BASE) pro = ret.dig[i]/BASE;
  ret.dig[i] -= BASE*pro;
 if(pro != 0) ret.dig.push_back(pro);
 return ret;
BigInt operator-(const BigInt& a) const {
 if(neg) return -(-(*this) - (-a));
  if(a.neg) return (*this) + (-a);
  int diff = cmp_minus(a);
  if(diff < 0) return -(a - (*this));</pre>
  if(diff == 0) return 0;
 BigInt ret; ret.dig.resize(len(), 0);
  for(int i=0;i<len();i++) {</pre>
  ret.dig[i] += dig[i];
```

```
if(i < a.len())    ret.dig[i] -= a.dig[i];
   if(ret.dig[i] < 0){</pre>
    ret.dig[i] += BASE;
    ret.dig[i+1]--;
  }
  ret.trim(); return ret;
 BigInt operator*(const BigInt& a) const {
  if(!len()||!a.len()) return 0;
  BigInt ret; ret.dig.resize(len()+a.len()+1);
  ret.neg = neg ^ a.neg;
  for(int i=0;i<len();i++)</pre>
   for(int j=0;j<a.len();j++){</pre>
    ret.dig[i+j] += dig[i] * a.dig[j];
    if(ret.dig[i+j] >= BASE) {
     lld x = ret.dig[i+j] / BASE;
     ret.dig[i+j+1] += x;
     ret.dig[i+j] -= x * BASE;
  ret.trim(); return ret;
 BigInt operator/(const BigInt& a) const {
  assert(a.len());
  if(len() < a.len()) return 0;</pre>
  BigInt ret; ret.dig.resize(len()-a.len()+1);
  ret.neg = a.neg;
  for(int i=len()-a.len();i>=0;i--){
   11d 1 = 0, r = BASE;
   while(r-1 > 1){
    11d \ mid = (1+r)>>1;
    ret.dig[i] = mid;
    if(ret*a<=(neg?-(*this):(*this))) 1 = mid;</pre>
    else r = mid;
   ret.dig[i] = 1;
  ret.neg ^= neg; ret.trim();
  return ret;
 BigInt operator%(const BigInt& a) const {
  return (*this) - (*this) / a * a;
 friend BigInt abs(BigInt a) { a.neg = 0; return a; }
friend void swap(BigInt& a, BigInt& b){
  swap(a.dig, b.dig); swap(a.neg, b.neg);
 friend istream& operator>>(istream& ss, BigInt& a){
  string s; ss >> s; a = s; return ss;
 friend ostream&operator<<(ostream&o, const BigInt&a){</pre>
  if(a.len() == 0) return o << '0';</pre>
  if(a.neg) o <<</pre>
  ss << o.dig.back()
  for(int i=a.len()-2;i>=0;i--)
   o<<setw(LOG_BASE)<<setfill('0')<<a.dig[i];
  return o;
 inline void print() const {
  if(len() == 0){putchar('0');return;}
  if(neg) putchar('-');
printf("%" PRINTF_ARG, dig.back());
  for(int i=len()-2;i>=0;i--)
printf("%0" LOG_BASE_STR PRINTF_ARG, dig[i]);
 #undef PRINTF_ARG
 #undef LOG_BASE_STR
}:
2.2 Dark Magic
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using __gnu_pbds::pairing_heap_tag;
using __gnu_pbds::binary_heap_tag;
using __gnu_pbds::binomial_heap_tag;
using __gnu_pbds::rc_binomial_heap_tag;
       __gnu_pbds::thin_heap_tag;
using
template<typename T>
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>,\
                     pairing_heap_tag>;
```

// a.join(b), pq.modify(pq.push(10), 87)

```
using __gnu_pbds::rb_tree_tag;
                                                             void to_child(Node* p,Node* c,bool dir){
using __gnu_pbds::ov_tree_tag;
                                                              p->ch[dir]=c;
using __gnu_pbds::splay_tree_tag;
                                                              p->up();
template<typename T>
using ordered_set = __gnu_pbds::tree<T,\</pre>
                                                             inline void rotate(Node* node){
__gnu_pbds::null_type,less<T>,rb_tree_tag,\
                                                              Node* par=node->par;
                                                              Node* par_par=par->par;
__gnu_pbds::tree_order_statistics_node_update>;
                                                              bool dir=node->is_rch();
// find_by_order, order_of_key
template<typename A, typename B>
                                                              bool par_dir=par->is_rch();
using hTable1=__gnu_pbds::cc_hash_table<A,B>;
                                                              to_child(par, node->ch[!dir], dir);
template<typename A, typename B>
                                                              to_child(node,par,!dir);
using hTable2=__gnu_pbds::gp_hash_table<A,B>;
                                                              if(par_par!=nullptr && par_par->ch[par_dir]==par)
                                                               to_child(par_par,node,par_dir);
2.3 Disjoint Set
                                                              else node->par=par_par;
class DJS {
                                                             inline void splay(Node* node){
private:
                                                              Node* tmp=node;
vector< int > fa, sz, sv;
vector< pair< int*, int > > opt;
                                                              stk[top++]=node;
void assign( int *k, int v ) {
                                                              while(!tmp->is_root()){
 opt.emplace_back( k, *k );
                                                               tmp=tmp->par;
                                                               stk[top++]=tmp;
  *k = v;
public:
                                                              while(top) stk[--top]->down();
                                                              for(Node *fa=node->par)
void init( int n ) {
  fa.resize( n ); iota( fa.begin(), fa.end(), 0 );
                                                               !node->is_root();
 sz.resize( n ); fill( sz.begin(), sz.end(), 1 );
                                                               rotate(node), fa=node->par)
                                                               if(!fa->is_root())
  opt.clear();
                                                                rotate(fa->is_rch()==node->is_rch()?fa:node);
int query(int x) {return fa[x] == x?x:query(fa[x]);}
void merge( int a, int b ) {
                                                             inline void access(Node* node){
 int af = query( a ), bf = query( b );
                                                              Node* last=nullptr;
  if( af == bf ) return;
                                                              while(node!=nullptr){
  if( sz[ af ] < sz[ bf ] ) swap( af, bf );</pre>
                                                               splay(node)
 assign( &fa[ bf ], fa[ af ] );
                                                               to_child(node, last, true);
 assign( &sz[ af ], sz[ af ] + sz[ bf ] );
                                                               last=node;
                                                               node=node->par;
void save() { sv.push_back( (int) opt.size() ); }
void undo() {
  int ls = sv.back(); sv.pop_back();
                                                             inline void change_root(Node* node){
 while ( ( int ) opt.size() > ls )
                                                              access(node);splay(node);node->set_rev();
  pair< int*, int > cur = opt.back();
   *cur.first = cur.second;
                                                             inline void link(Node* x, Node* y){
   opt.pop_back();
                                                              change_root(x);splay(x);x->par=y;
 }
}
                                                             inline void split(Node* x, Node* y) {
                                                              {\tt change\_root(x);access(y);splay(x)}
};
                                                              to_child(x,nullptr,true);y->par=nullptr;
     Link-Cut Tree
struct Node{
                                                             inline void change_val(Node* node,int v){
Node *par, *ch[2];
                                                              access(node);splay(node);node->v=v;node->up();
 int xor_sum, v;
 bool is_rev;
                                                             inline int query(Node* x,Node* y){
                                                              change\_root(x); access(y); splay(y);
Node(int _v){
  v=xor_sum=_v;is_rev=false;
                                                              return y->xor_sum;
 par=ch[0]=ch[1]=nullptr;
                                                             inline Node* find_root(Node* node){
inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
                                                              access(node);splay(node);
inline void down(){
                                                              Node* last=nullptr;
 if(is_rev){
                                                              while(node!=nullptr){
   if(ch[0]!=nullptr) ch[0]->set_rev();
                                                               node->down();last=node;node=node->ch[0];
   if(ch[1]!=nullptr) ch[1]->set_rev();
   is_rev=false;
                                                              return last;
 }
                                                             set<pii> dic;
 inline void up(){
                                                             inline void add_edge(int u,int v){
 xor_sum=v;
                                                              if(u>v) swap(u,v)
                                                              if(find_root(node[u])==find_root(node[v])) return;
  if(ch[0]!=nullptr){
  xor_sum^=ch[0]->xor_sum;
                                                              dic.insert(pii(u,v))
  ch[0]->par=this;
                                                              link(node[u],node[v]);
                                                             inline void del_edge(int u,int v){
  if(ch[1]!=nullptr){
                                                              if(u>v) swap(u,v);
  xor_sum^=ch[1]->xor_sum;
   ch[1]->par=this;
                                                              if(dic.find(pii(u,v))==dic.end()) return;
  }
                                                              dic.erase(pii(u,v))
                                                              split(node[u],node[v]);
inline bool is_root(){
  {\color{red} \textbf{return}} \  \, {\color{blue} \textbf{par} = = \textbf{nullptr}} \  \, |\,|\,\backslash \\
                                                             2.5 LiChao Segment Tree
   (par->ch[0]!=this && par->ch[1]!=this);
                                                             struct Line{
bool is_rch(){return !is_root() && par->ch[1]==this;}
                                                              int m, k, id;
} *node[maxn], *stk[maxn];
                                                              Line() : id( -1 ) {}
int top;
                                                              Line( int a, int b, int c )
```

```
: m( a ), k( b ), id( c ) {}
                                                             template < typename T, typename Cmp_ = less< T > >
 int at( int x ) { return m * x + k; }
                                                             class SparseTable {
                                                             private:
class LiChao {
                                                              vector< vector< T > > tbl;
                                                              vector< int > lg;
private:
 int n; vector< Line > nodes;
                                                              T cv( T a, T b ) {
  inline int lc( int x ) { return 2 * x + 1; }
                                                               return Cmp_()( a, b ) ? a : b;
  inline int rc( int x ) { return 2 * x + 2; }
 void insert( int 1, int r, int id, Line ln ) {
  int m = (1 + r) >> 1;
                                                             public:
                                                              void init( T arr[], int n ) {
   if ( nodes[ id ].id == -1 ) {
                                                               // 0-base
   nodes[ id ] = ln;
                                                               lg.resize(n+1);
                                                               lg[0] = -1;
    return:
                                                               for( int i=1 ; i<=n ; ++i ) lg[i] = lg[i>>1] + 1;
                                                               tbl.resize(lg[n] + 1);
   bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
   if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
                                                               tbl[ 0 ].resize( n )
   atLeft ^= 1; swap( nodes[ id ], ln );
                                                               copy( arr, arr + n, tbl[ 0 ].begin() );
                                                               for ( int i = 1 ; i <= lg[ n ] ; ++ i ) {
   if ( r - 1 == 1 ) return;
                                                                int len = 1 << ( i - 1 ), sz = 1 << i;
                                                                tbl[ i ].resize( n - sz + 1 );
for ( int j = 0 ; j <= n - sz ; ++ j
   if ( atLeft ) insert( 1, m, lc( id ), ln );
   else insert( m, r, rc( id ), ln );
                                                                 tbl[i][j] = cv(tbl[i-1][j], tbl[i-1][j+len]);
  int query( int 1, int r, int id, int x ) {
  int ret = 0;
   if ( nodes[ id ].id != -1 )
                                                              T query( int 1, int r ) {
                                                               // 0-base [1, r)
    ret = nodes[ id ].at( x );
                                                               int wh = lg[ r - l ], len = 1 << wh;</pre>
   int m = (1 + r) >> 1;
   if ( r - l == 1 ) return ret;
                                                               return cv( tbl[ wh ][ 1 ], tbl[ wh ][ r - len ] );
   else if ( x < m )</pre>
                                                             };
    return max( ret, query( 1, m, lc( id ), x ) );
    return max( ret, query( m, r, rc( id ), x ) );
                                                                   Linear Basis
                                                             2.8
public:
                                                             struct LinearBasis {
 void build( int n_ ) {
                                                             private:
  n = n_; nodes.clear();
                                                              int n, sz;
  nodes.resize( n << 2, Line() );</pre>
                                                              vector< llu > B;
                                                              inline llu two( int x ){ return ( ( llu ) 1 ) << x; }</pre>
  void insert( Line ln ) { insert( 0, n, 0, ln ); }
                                                             public:
 int query( int x ) { return query( 0, n, 0, x ); }
                                                              void init( int n_ ) {
} lichao;
                                                               n = n_{;} B.clear(); B.resize(n); sz = 0;
2.6 Treap
                                                              void insert( llu x ) {
namespace Treap{
                                                               // add x into B
#define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                               for ( int i = n-1; i >= 0; --i ) if( two(i) & x ){
struct node{
                                                                if ( B[ i ] ) x ^= B[ i ];
  int size;
                                                                else {
 uint32_t pri;
                                                                 B[ i ] = x; sz++;
 node *lc, *rc;
                                                                 for ( int j = i - 1 ; j >= 0 ; -- j )
if( B[ j ] && ( two( j ) & B[ i ] ))
 node() : size(0), pri(rand()), lc(0), rc(0) {}
 void pull() {
                                                                    B[ i ] ^= B[ j ];
  size = 1;
                                                                 for (int j = i + 1; j < n; ++ j)
if (two(i) & B[j])
   if ( lc ) size += lc->size;
  if ( rc ) size += rc->size;
                                                                   B[ j ] ^= B[ i ];
 }
                                                                 break;
 }:
node* merge( node* L, node* R ) {
                                                               }
 if ( not L or not R ) return L ? L : R;
 if ( L->pri > R->pri ) {
                                                              inline int size() { return sz; }
  L->rc = merge( L->rc, R ); L->pull();
                                                              bool check( llu x )
  return L;
                                                               // is x in span(B) ?
  } else {
                                                               for ( int i = n-1 ; i >= 0 ; --i ) if( two(i) & x )
if( B[ i ] ) x ^= B[ i ];
  R->lc = merge( L, R->lc ); R->pull();
   return R:
                                                                else return false;
 }
                                                               return true:
void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                              llu kth_small(llu k) {
 if ( not rt ) L = R = nullptr;
                                                               /** 1-base would always > 0 **/
  else if( sz( rt->lc ) + 1 <= k ) {
                                                                /** should check it **/
                                                               /* if we choose at least one element
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                 but size(B)(vectors in B)==N(original elements)
   L->pull();
                                                                  then we can't get 0 */
  } else {
                                                               11u ret = 0;
  R = rt;
                                                               for ( int i = 0 ; i < n ; ++ i ) if( B[ i ] ) {
   split_by_size( rt->lc, k, L, R->lc );
                                                                if( k & 1 ) ret ^= B[ i ];
   R->pull();
                                                                k >>= 1;
 }
                                                               }
                                                               return ret;
 #undef sz
```

} base;

2.7 Sparse Table

3 Graph

3.1 Euler Circuit

```
bool vis[ N ]; size_t la[ K ];
void dfs( int u, vector< int >& vec ) {
while ( la[ u ] < G[ u ].size() ) {</pre>
  if( vis[ G[ u ][ la[ u ] ].second ] ) {
   ++ la[ u ];
   continue;
 int v = G[ u ][ la[ u ] ].first;
vis[ G[ u ][ la[ u ] ].second ] = true;
++ la[ u ]; dfs( v, vec );
  vec.push_back( v );
3.2 BCC Edge
class BCC{
private:
 vector< int > low, dfn;
 int cnt:
 vector< bool > bridge;
 vector< vector< PII > > G;
 void dfs( int w, int f ) {
  low[ w ] = dfn[ w ] = cnt++
  for ( auto [ u, t ] : G[ w ] ) {
   if ( u == f ) continue;
   if ( dfn[ u ] != 0 ) {
  low[ w ] = min( low[ w ], dfn[ u ] );
   }else{
    dfs( u, w );
    low[ w ] = min( low[ w ], low[ u ] );
if ( low[ u ] > dfn[ w ] ) bridge[ t ] = true;
  }
public:
 void init( int n, int m ) {
  G.resize(n); cnt = 0;
  fill( G.begin(), G.end(), vector< PII >() );
  bridge.clear(); bridge.resize( m );
  low.clear(); low.resize( n );
  dfn.clear(); dfn.resize( n );
 void add_edge( int u, int v ) {
  // should check for multiple edge
  G[ u ].emplace_back( v, cnt );
  G[ v ].emplace_back( u, cnt ++ );
 }
 void solve(){
  cnt = 1;
  for (int i = 0; i < n; ++i)</pre>
   if (not vis[ i ]) dfs(i, i);
 // the id will be same as insert order, \theta\text{-base}
 bool is_bridge( int x ) { return bridge[ x ]; }
3.3 BCC Vertex
class BCC {
  private:
    int n, t, ecnt;
    vector<vector<pair<int, int>>> G;
    vector<int> low, tin, st, bcc;
    vector<bool> ap, ins;
void dfs(int x, int p)
       tin[x] = low[x] = ++t;
       int ch = 0;
      for (auto u: G[x]) {
  if (u.first == p) continue;
         if (not ins[u.second]) {
           st.push_back(u.second);
           ins[u.second] = true;
         if (tin[u.first])
           low[x] = min(low[x], tin[u.first]);
           continue;
         ++ch; dfs(u.first, x);
         low[x] = min(low[x], low[u.first]);
```

```
if (low[u.first] >= tin[x]) {
          ap[x] = true; ++ecnt;
          while (true) {
            int e = st.back(); st.pop_back();
            bcc[e] = ecnt;
            if (e == u.second) break;
        }
      if (ch == 1 \text{ and } p == x) ap[x] = false;
  public:
    void init(int n_) {
      n = n_, ecnt = 0; st.clear();
      G.clear(); G.resize(n);
      low.clear(); tin.clear();
      ap.assign(n, false);
    void add_edge(int u, int v) {
      G[u].emplace_back(v, ecnt);
      G[v].emplace_back(u, ecnt++);
    void solve() {
      ecnt = 0; bcc.resize(t);
      ins.assign(t, false);
      for (int i = 0; i < n; ++i)
        if (low[i] == 0) dfs(i, i);
    int get_id(int x) { return bcc[x];; }
    int count() { return ecnt; }
    bool is_ap(int x) { return ap[x]; }
};
3.4 2-SAT (SCC)
class TwoSat{
 private:
  int n:
  vector<vector<int>> rG,G,sccs;
  vector<int> ord,idx;
  vector<bool> vis,result;
  void dfs(int u){
   vis[u]=true
   for(int v:G[u])
   if(!vis[v]) dfs(v);
   ord.push_back(u);
  void rdfs(int u){
   vis[u]=false;idx[u]=sccs.size()-1;
   sccs.back().push_back(u);
   for(int v:rG[u])
    if(vis[v])rdfs(v);
 public:
  void init(int n_){
   n=n_;G.clear();G.resize(n);
   rG.clear();rG.resize(n)
   sccs.clear();ord.clear();
   idx.resize(n);result.resize(n);
  void add_edge(int u,int v){
   G[u].push_back(v);rG[v].push_back(u);
  void orr(int x,int y){
   if ((x^y)==1)return
   add_edge(x^1,y); add_edge(y^1,x);
  bool solve(){
   vis.clear();vis.resize(n);
   for(int i=0;i<n;++i)</pre>
    if(not vis[i])dfs(i);
   reverse(ord.begin(),ord.end());
   for (int u:ord){
    if(!vis[u])continue;
    sccs.push_back(vector<int>());
    rdfs(u);
   for(int i=0;i<n;i+=2)</pre>
    if(idx[i]==idx[i+1])
     return false;
   vector<bool> c(sccs.size());
   for(size_t i=0;i<sccs.size();++i){</pre>
    for(size_t j=0;j<sccs[i].size();++j){</pre>
```

```
result[sccs[i][j]]=c[i];
                                                                     dfschain(1,1);
     c[idx[sccs[i][j]^1]]=!c[i];
                                                                    PII get_inter( int u ) { return {tl[ u ], tr[ u ]}; }
                                                                    vector< PII > get_path( int u , int v ){
                                                                     vector< PII > res;
   return true;
                                                                     int g = lca( u, v );
while ( chain[ u ] != chain[ g ] ) {
  bool get(int x){return result[x];}
                                                                      int s = chain_st[ chain[ u ] ];
  inline int get_id(int x){return idx[x];}
  inline int count(){return sccs.size();}
                                                                      res.emplace_back( tl[ s ], tl[ u ] + 1 );
                                                                      u = fa[ s ][ 0 ];
} sat2:
                                                                     res.emplace_back( tl[ g ], tl[ u ] + 1 );
while ( chain[ v ] != chain[ g ] ) {
  int s = chain_st[ chain[ v ] ];
3.5 Lowbit Decomposition
class LowbitDecomp{
private:
int time_, chain_, LOG_N;
                                                                      res.emplace_back( tl[ s ], tl[ v ] + 1 );
vector< vector< int > > G, fa;
                                                                      v = fa[ s ][ 0 ];
vector< int > tl, tr, chain, chain_st;
// chain_ : number of chain
                                                                     res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
// tl, tr[ u ] : subtree interval in the seq. of u
                                                                     return res;
// chain_st[ u ] : head of the chain contains u
// chian[ u ] : chain id of the chain u is on
                                                                     /* res : list of intervals from u to v
                                                                      * ( note only nodes work, not edge )
                                                                      * usage
inline int lowbit( int x ) {
  return x & ( -x );
                                                                      * vector< PII >& path = tree.get_path( u , v )
                                                                      * for( auto [ 1, r ] : path ) {
                                                                      * 0-base [ 1, r )
void predfs( int u, int f ) {
  chain[ u ] = 0;
                                                                      */
  for ( int v : G[ u ] ) {
  if ( v == f ) continue;
   predfs( v, u );
                                                                  } tree;
   if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
                                                                        MaxClique
    chain[ u ] = chain[ v ];
                                                                   // contain a self loop u to u, than u won't in clique
  if ( not chain[ u ] )
                                                                   template < size_t MAXN >
                                                                  class MaxClique{
   chain[ u ] = chain_ ++;
                                                                   private:
 void dfschain( int u, int f ) {
                                                                    using bits = bitset< MAXN >;
  fa[ u ][ 0 ] = f;
                                                                    bits popped, G[ MAXN ], ans;
                                                                    size_t deg[ MAXN ], deo[ MAXN ], n;
  for ( int i = 1 ; i < LOG_N ; ++ i )
   fa[u][i] = fa[fa[u][i-1]][i-1];
                                                                    void sort_by_degree() {
  tl[ u ] = time_++
                                                                     popped.reset();
  if ( not chain_st[ chain[ u ] ] )
                                                                     for ( size_t i = 0 ; i < n ; ++ i )</pre>
  chain_st[ chain[ u ] ] = u;
                                                                       deg[ i ] = G[ i ].count();
                                                                     for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
  for ( int v : G[ u ] )
  if ( v != f and chain[ v ] == chain[ u ] )
                                                                       size_t mi = MAXN, id = 0;
                                                                       for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )</pre>
  dfschain( v, u );
for ( int v : G[ u ] )
   if ( v != f and chain[ v ] != chain[ u ] )
                                                                           mi = deg[id = j];
                                                                       popped[ deo[ i ] = id ] = 1;
for( size_t u = G[ i ]._Find_first() ;
    dfschain( v, u );
  tr[ u ] = time_;
                                                                        u < n ; u = G[ i ]._Find_next( u ) )</pre>
inline bool anc( int u, int v ) {
  return tl[ u ] <= tl[ v ] \
  and tr[ v ] <= tr[ u ];</pre>
                                                                          -- deg[ u ];
                                                                     }
                                                                    void BK( bits R, bits P, bits X ) {
                                                                     if (R.count()+P.count() <= ans.count()) return;</pre>
public:
inline int lca( int u, int v ) {
                                                                     if ( not P.count() and not X.count() )
  if ( anc( u, v ) ) return u;
                                                                      if ( R.count() > ans.count() ) ans = R;
  for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
                                                                      return:
   if ( not anc( fa[ u ][ i ], v ) )
    u = fa[ u ][ i ];
                                                                     /* greedily chosse max degree as pivot
                                                                     bits cur = P | X; size_t pivot = 0, sz = 0;
for ( size_t u = cur._Find_first() ;
  return fa[ u ][ 0 ];
void init( int n ) {
                                                                      u < n ; u = cur._Find_next( u ) )</pre>
  \begin{array}{l} n \ ++; \\ \text{for} \ ( \ \text{LOG\_N} \ = \ 0 \ ; \ ( \ 1 \ << \ \text{LOG\_N} \ ) \ < \ n \ ; \ ++ \ \text{LOG\_N} \ ); \\ \end{array} 
                                                                       if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
                                                                     cur = P & ( ~G[ pivot ] );
  fa.clear();
                                                                     */ // or simply choose first
                                                                     bits cur = P & (~G[ ( P | X )._Find_first() ]);
  fa.resize( n, vector< int >( LOG_N ) );
 G.clear(); G.resize( n );
tl.clear(); tl.resize( n );
                                                                     for ( size_t u = cur._Find_first()
                                                                      u < n ; u = cur._Find_next( u ) ) {
                                                                      if ( R[ u ] ) continue;
  tr.clear(); tr.resize( n );
  chain.clear(); chain.resize( n );
                                                                      R[u] = 1;
                                                                      BK( R, P & G[ u ], X & G[ u ] );
  chain_st.clear(); chain_st.resize( n );
                                                                      R[u] = P[u] = 0, X[u] = 1;
void add_edge( int u , int v ) {
  // 1-base
  G[ u ].push_back( v );
                                                                  public:
 G[ v ].push_back( u );
                                                                    void init( size_t n_ ) {
                                                                     n = n_{-};
void decompose(){
                                                                     for ( size_t i = 0 ; i < n ; ++ i )
  chain_ = 1;
                                                                      G[ i ].reset();
  predfs( 1, 1 );
                                                                     ans.reset();
  time_{-} = 0;
```

```
void add_edges( int u, bits S ) { G[ u ] = S; }
                                                                    sort(r.begin(), r.end(),
                                                                     [&](int i, int j) { return d[i] > d[j]; });
void add_edge( int u, int v ) {
 G[u][v] = G[v][u] = 1;
                                                                    csort(r, c);
                                                                    dfs(r, c, 1, mask);
return ans; // sol[0 ~ ans-1]
int solve() {
 sort_by_degree(); // or simply iota( deo... )
for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                 } graph;
   deg[ i ] = G[ i ].count();
                                                                  3.8 Virtural Tree
 bits pob, nob = 0; pob.set();
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
                                                                 inline bool cmp(const int &i, const int &j) {
 for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
                                                                   return dfn[i] < dfn[j];</pre>
   size_t v = deo[ i ];
  bits tmp; tmp[ v ] = 1;
BK( tmp, pob & G[ v ], nob & G[ v ] );
pob[ v ] = 0, nob[ v ] = 1;
                                                                 void build(int vectrices[], int k) {
                                                                   static int stk[MAX_N];
                                                                   sort(vectrices, vectrices + k, cmp);
                                                                   stk[sz++] = 0;
  return static_cast< int >( ans.count() );
                                                                   for (int i = 0; i < k; ++i) {
  int u = vectrices[i], lca = LCA(u, stk[sz - 1]);
  if (1)</pre>
                                                                    if (lca == stk[sz - 1]) stk[sz++] = u;
     MaxCliqueDyn
3.7
                                                                     while (sz \ge 2 \& dep[stk[sz - 2]] \ge dep[lca]) {
constexpr int kN = 150;
                                                                      addEdge(stk[sz - 2], stk[sz - 1]);
struct MaxClique { // Maximum Clique
bitset<kN> a[kN], cs[kN];
                                                                      sz--:
int ans, sol[kN], q, cur[kN], d[kN], n;
                                                                     if (stk[sz - 1] != lca) {
void init(int _n) {
                                                                      addEdge(lca, stk[--sz]);
 n = _n; for (int i = 0; i < n; i++) a[i].reset();
                                                                      stk[sz++] = lca, vectrices[cnt++] = lca;
void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
void csort(vector<int> &r, vector<int> &c) {
                                                                     stk[sz++] = u;
 int mx = 1, km = max(ans - q + 1, 1), t = 0,
    m = int(r.size())
                                                                   for (int i = 0; i < sz - 1; ++i)
  cs[1].reset(); cs[2].reset()
                                                                    addEdge(stk[i], stk[i + 1]);
 for (int i = 0; i < m; i++) {
   int p = r[i], k = 1
   while ((cs[k] & a[p]).count()) k++;
                                                                  3.9 Tree Hashing
   if (k > mx) cs[++mx + 1].reset();
                                                                 uint64_t hsah( int u, int f ) {
   cs[k][p] = 1;
                                                                    uint64_t r = 127;
   if (k < km) r[t++] = p;
                                                                    for ( int v : G[ u ] ) {
  if ( v == f ) continue;
 c.resize(m);
                                                                      uint64_t hh = hsah( v, u );
 if(t) c[t-1] = 0;
                                                                      r = r + (hh * hh) % mod;
  for (int k = km; k <= mx; k++) {</pre>
  for (int p = int(cs[k]._Find_first());
                                                                    return r:
      p < kN; p = int(cs[k]._Find_next(p))) {
                                                                 }
    r[t] = p; c[t++] = k;
                                                                         Minimum Mean Cycle
  }
                                                                  /* minimum mean cycle O(VE) */
void dfs(vector<int> &r, vector<int> &c, int 1,
                                                                 struct MMC{
 bitset<kN> mask) {
                                                                  #define FZ(n) memset((n),0,sizeof(n))
 while (!r.empty()) {
                                                                  #define E 101010
   int p = r.back(); r.pop_back();
                                                                  #define V 1021
   mask[p] = 0;
                                                                 #define inf 1e9
   if (q + c.back() <= ans) return;</pre>
                                                                   struct Edge { int v,u; double c; };
   cur[q++] = p;
                                                                   int n, m, prv[V][V], prve[V][V], vst[V];
   vector<int> nr, nc;
                                                                   Edge e[E];
                                                                   vector<int> edgeID, cycle, rho;
   bitset<kN> nmask = mask & a[p];
   for (int i : r)
                                                                   double d[V][V];
    if (a[p][i]) nr.push_back(i);
                                                                   void init( int _n ) { n = _n; m = 0; }
                                                                   // WARNING: TYPE matters
   if (!nr.empty()) {
    if (1 < 4) {
                                                                   void add_edge( int vi , int ui , double ci )
     for (int i : nr)
                                                                   { e[ m ++ ] = { vi , ui , ci }; }
      d[i] = int((a[i] & nmask).count());
                                                                   void bellman_ford() {
                                                                    for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {</pre>
     sort(nr.begin(), nr.end(),
      [&](int x, int y)
                                                                     fill(d[i+1], d[i+1]+n, inf);
for(int j=0; j<m; j++) {
  int v = e[j].v, u = e[j].u;</pre>
        return d[x] > d[y];
      });
   csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
                                                                      if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
                                                                       d[i+1][u] = d[i][v]+e[j].c;
                                                                       prv[i+1][u] = v
    ans = q; copy(cur, cur + q, sol);
                                                                       prve[i+1][u] = j;
   c.pop_back(); q--;
  }
int solve(bitset<kN> mask) { // vertex mask
 vector<int> r, c;
                                                                   double solve(){
  for (int i = 0; i < n; i++)
                                                                    // returns inf if no cycle, mmc otherwise
  if (mask[i]) r.push_back(i);
for (int i = 0; i < n; i++)</pre>
                                                                    double mmc=inf;
                                                                    int st = -1;
   d[i] = int((a[i] & mask).count());
                                                                    bellman_ford();
```

```
for(int i=0; i<n; i++) {</pre>
                                                                   // Minimum Steiner Tree
                                                                   // O(V 3^T + V^2 2^T)
   double avg=-inf;
   for(int k=0; k<n; k++) {</pre>
                                                                   struct SteinerTree{
    if(d[n][i]<inf-eps)</pre>
                                                                   #define V 33
     avg=max(avg,(d[n][i]-d[k][i])/(n-k));\\
                                                                   #define T 8
                                                                   #define INF 1023456789
    else avg=max(avg,inf);
                                                                     int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
                                                                     void init( int _n ){
                                                                      n = _n;
                                                                      for( int i = 0 ; i < n ; i ++ ){
  for( int j = 0 ; j < n ; j ++ )</pre>
  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
  for (int i=n; !vst[st]; st=prv[i--][st]) {
   vst[st]++;
                                                                        dst[ i ][ j ] = INF;
   edgeID.PB(prve[i][st]);
                                                                       dst[ i ][ i ] = 0;
   rho.PB(st);
                                                                    void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
  dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
  while (vst[st] != 2) {
   int v = rho.back(); rho.pop_back();
   cycle.PB(v);
   vst[v]++;
                                                                     void shortest_path(){
                                                                      for( int k = 0 ; k < n ; k ++ )
 reverse(ALL(edgeID));
                                                                       for( int i = 0 ; i < n ; i ++ )</pre>
  edgeID.resize(SZ(cycle));
                                                                        for( int j = 0 ; j < n ; j ++ )
dst[ i ][ j ] = min( dst[ i ][ j ],</pre>
  return mmc;
 }
} mmc;
                                                                             dst[ i ][ k ] + dst[ k ][ j ] );
3.11 Mo's Algorithm on Tree
                                                                     int solve( const vector<int>& ter ){
                                                                      int t = (int)ter.size();
int q; vector< int > G[N];
                                                                     for( int i = 0 ; i < (1 << t ) ; i ++ )
for( int j = 0 ; j < n ; j ++ )
dp[ i ][ j ] = INF;
struct Que{
int u, v, id;
} que[ N ];
                                                                      for( int i = 0 ; i < n ; i ++ )</pre>
int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
                                                                       dp[0][i] = 0;
void dfs( int u, int f ) {
                                                                      for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
dfn[ u ] = dfn_++; int saved_rbp = stk_;
for ( int v : G[ u ] ) {
                                                                       if( msk == ( msk & (-msk) ) ){
                                                                        int who = __lg( msk );
for( int i = 0 ; i < n ; i ++ )</pre>
  if ( v == f ) continue;
  dfs( v, u );
                                                                         dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
  if ( stk_ - saved_rbp < SQRT_N ) continue;</pre>
                                                                        continue:
  for ( ++ block_ ; stk_ != saved_rbp ; )
  block_id[ stk[ -- stk_ ] ] = block_;
                                                                       for( int i = 0 ; i < n ; i ++ )</pre>
                                                                        for( int submsk = ( msk - 1 ) & msk ; submsk ;
stk[ stk_ ++ ] = u;
                                                                              submsk = ( submsk - 1 ) & msk )
                                                                           dp[ msk ][ i ] = min( dp[ msk ][ i ],
bool inPath[ N ];
                                                                                    dp[ submsk ][ i ] +
void Diff( int u ) {
                                                                                    dp[ msk ^ submsk ][ i ] );
if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
                                                                       for( int i = 0 ; i < n ; i ++ ){</pre>
 else { /*add this edge*/ }
                                                                        tdst[ i ] = INF;
for( int j = 0 ; j < n ; j ++ )
tdst[ i ] = min( tdst[ i ],</pre>
void traverse( int& origin_u, int u ) {
for ( int g = lca( origin_u, u )
                                                                                dp[ msk ][ j ] + dst[ j ][ i ] );
 origin_u != g ; origin_u = parent_of[ origin_u ] )
   Diff( origin_u );
                                                                       for( int i = 0 ; i < n ; i ++ )</pre>
 for (int v = u; v != origin_u; v = parent_of[v])
                                                                        dp[ msk ][ i ] = tdst[ i ];
 Diff( v );
 origin_u = u;
                                                                      int ans = INF:
                                                                      for( int i = 0 ; i < n ; i ++ )</pre>
void solve() {
                                                                       ans = min(ans, dp[(1 << t) - 1][i]);
 dfs( 1, 1 );
                                                                      return ans:
 while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
 sort( que, que + q, [](const Que& x, const Que& y) {
                                                                   } solver;
  return tie( block_id[ x.u ], dfn[ x.v ] )
       < tie( block_id[ y.u ], dfn[ y.v ] );
                                                                    3.13 Directed Minimum Spanning Tree
 } );
                                                                    template <typename T> struct DMST {
 int U = 1, V = 1;
 for ( int i = 0 ; i < q ; ++ i ) {
  pass( U, que[ i ].u );</pre>
                                                                     T g[maxn][maxn], fw[maxn];
                                                                     int n, fr[maxn];
  pass( V, que[ i ].v );
                                                                     bool vis[maxn], inc[maxn];
                                                                     void clear() {
  // we could get our answer of que[ i ].id
                                                                      for(int i = 0; i < maxn; ++i) {</pre>
                                                                       for(int j = 0; j < maxn; ++j) g[i][j] = inf;
}
                                                                       vis[i] = inc[i] = false;
Method 2:
dfs u:
                                                                     void addEdge(int u,int v,T w){g[u][v]=min(g[u][v],w);}
push u
                                                                     T operator()(int root, int _n) {
 iterate subtree
                                                                     n = _n; T ans = 0;
 push u
                                                                      if (dfs(root) != n) return -1;
Let P = LCA(u, v), and St(u) <= St(v)
if (P == u) query[St(u), St(v)]
                                                                      while (true) {
                                                                       for(int i = 1;i <= n;++i) fw[i] = inf, fr[i] = i;</pre>
else query[Ed(u), St(v)], query[St(P), St(P)]
                                                                       for (int i = 1; i <= n; ++i) if (!inc[i]) {
  for (int j = 1; j <= n; ++j) {</pre>
                                                                         if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
```

3.12 Minimum Steiner Tree

```
fw[i] = g[j][i]; fr[i] = j;
                                                                 if (sdom[p] == i) dom[u] = i;
                                                                else dom[u] = p;
    }
                                                               }
                                                               if (i) merge(i, rp[i]);
   int x = -1:
   for(int i = 1;i <= n;++i)if(i != root && !inc[i]){</pre>
                                                              vector<int> p(n, -2); p[s] = -1;
                                                              for (int i = 1; i < tk; ++i)
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];</pre>
    int j = i, c = 0;
    while(j!=root && fr[j]!=i && c<=n) ++c, j=fr[j];</pre>
    if (j == root || c > n) continue;
                                                              for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
    else { x = i; break; }
                                                              return p:
                                                             }}
   if (!~x) {
    for (int i = 1; i <= n; ++i)</pre>
                                                             4
                                                                   Matching & Flow
     if (i != root && !inc[i]) ans += fw[i];
    return ans;
                                                             4.1
                                                                  Kuhn Munkres
                                                             class KM {
   int y = x;
                                                             private:
   for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
                                                              static constexpr lld INF = 1LL << 60;</pre>
   do {
                                                              vector<lld> hl,hr,slk;
    ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true;
                                                              vector<int> fl,fr,pre,qu;
   } while (y != x);
                                                              vector<vector<lld>> w;
   inc[x] = false;
                                                              vector<bool> v1,vr;
   for (int k = 1; k <= n; ++k) if (vis[k]) {
                                                              int n, ql, qr;
    for (int j = 1; j <= n; ++j) if (!vis[j]) {</pre>
                                                              bool check(int x) {
     if (g[x][j] > g[k][j]) g[x][j] = g[k][j]
                                                               if (vl[x] = true, fl[x] != -1)
     if (g[j][k] < \inf \&\& g[j][k]-fw[k] < g[j][x])
                                                                return vr[qu[qr++] = f1[x]] = true;
      g[j][x] = g[j][k] - fw[k];
                                                               while (x != -1) swap(x, fr[fl[x] = pre[x]]);
                                                               return false:
   }
                                                              void bfs(int s) {
  return ans;
                                                               fill(slk.begin(), slk.end(), INF);
                                                               fill(vl.begin(), vl.end(), false);
 int dfs(int now) {
                                                               fill(vr.begin(), vr.end(), false);
  int r = 1; vis[now] = true;
                                                               ql = qr = 0;
  for (int i = 1; i <= n; ++i)</pre>
                                                               qu[qr++] = s;
   if (g[now][i] < inf && !vis[i]) r += dfs(i);</pre>
                                                               vr[s] = true;
  return r;
                                                               while (true) {
                                                                11d d;
};
                                                                while (ql < qr) {
  for (int x = 0, y = qu[ql++]; x < n; ++x) {</pre>
3.14
      Dominator Tree
                                                                  if(!v1[x]&&slk[x]>=(d=h1[x]+hr[y]-w[x][y])){
namespace dominator {
                                                                   if (pre[x] = y, d) slk[x] = d;
vector<int> g[maxn], r[maxn], rdom[maxn];
                                                                    else if (!check(x)) return;
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
                                                                 }
void init(int n) {
 // vertices are numbered from 0 to n - 1
                                                                d = INF;
 fill(dfn, dfn + n, -1);fill(rev, rev + n, -1);
                                                                for (int x = 0; x < n; ++x)
 fill(fa, fa + n, -1); fill(val, val + n, -1);
                                                                  if (!vl[x] && d > slk[x]) d = slk[x];
 fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
                                                                 for (int x = 0; x < n; ++x) {
 fill(dom, dom + n, -1); tk = 0;
                                                                 if (v1[x]) h1[x] += d;
 for (int i = 0; i < n; ++i) {
                                                                  else slk[x] -= d;
  g[i].clear(); r[i].clear(); rdom[i].clear();
                                                                 if (vr[x]) hr[x] -= d;
                                                                for (int x = 0; x < n; ++x)
void add_edge(int x, int y) { g[x].push_back(y); }
                                                                 if (!v1[x] && !s1k[x] && !check(x)) return;
void dfs(int x) {
 rev[dfn[x] = tk] = x;
 fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
                                                             public:
for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
                                                              void init( int n_ ) {
                                                               n = n_; qu.resize(n);
  r[dfn[u]].push_back(dfn[x]);
                                                               fl.clear(); fl.resize(n, -1);
                                                               fr.clear(); fr.resize(n, -1);
                                                               hr.clear(); hr.resize(n); hl.resize(n);
void merge(int x, int y) { fa[x] = y; }
                                                               w.clear(); w.resize(n, vector<lld>(n));
int find(int x, int c = 0) {
                                                               slk.resize(n); pre.resize(n);
 if (fa[x] == x) return c ? -1 : x;
                                                               vl.resize(n); vr.resize(n);
int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
                                                              void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
                                                              1ld solve() {
 fa[x] = p;
                                                               for (int i = 0; i < n; ++i)
 return c ? p : val[x];
                                                                hl[i] = *max_element(w[i].begin(), w[i].end());
                                                               for (int i = 0; i < n; ++i) bfs(i);</pre>
vector<int> build(int s, int n) {
                                                               11d res = 0;
// return the father of each node in the dominator tree
                                                               for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
// p[i] = -2 if i is unreachable from s
                                                               return res;
 dfs(s);
 for (int i = tk - 1; i >= 0; --i)
                                                             } km;
  for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
  if (i) rdom[sdom[i]].push_back(i);
                                                             4.2 Bipartite Matching
  for (int &u : rdom[i]) {
   int p = find(u);
                                                            class BipartiteMatching{
```

```
private:
vector<int> X[N], Y[N];
int fX[N], fY[N], n;
bitset<N> walked;
bool dfs(int x){
  for(auto i:X[x]){
   if(walked[i])continue;
   walked[i]=1;
   if(fY[i]==-1||dfs(fY[i])){
    fY[i]=x;fX[x]=i;
    return 1:
  }
  return 0;
public:
void init(int _n){
 n=_n; walked.reset();
  for(int i=0;i<n;i++)</pre>
   X[i].clear();Y[i].clear();
   fX[i]=fY[i]=-1;
  }
void add_edge(int x, int y){
 X[x].push_back(y); Y[y].push_back(y);
int solve(){
 int cnt = 0;
  for(int i=0;i<n;i++){</pre>
  walked.reset();
   if(dfs(i)) cnt++;
  // return how many pair matched
  return cnt;
};
      General Graph Matching
const int N = 514, E = (2e5) * 2;
struct Graph{
int to[E],bro[E],head[N],e;
int lnk[N], vis[N], stp,n;
void init( int _n ){
 stp = 0; e = 1; n = _n;
  for( int i = 0 ; i <= n ; i ++ )</pre>
  head[i] = lnk[i] = vis[i] = 0;
void add_edge(int u,int v){
  to[e]=v,bro[e]=head[u],head[u]=e++;
 to[e]=u,bro[e]=head[v],head[v]=e++;
bool dfs(int x){
  vis[x]=stp;
  for(int i=head[x];i;i=bro[i]){
  int v=to[i];
   if(!lnk[v]){
    lnk[x]=v, lnk[v]=x;
    return true
   }else if(vis[lnk[v]]<stp){</pre>
    int w=lnk[v];
    lnk[x]=v, lnk[v]=x, lnk[w]=0;
    if(dfs(w)) return true
    lnk[w]=v, lnk[v]=w, lnk[x]=0;
 }
  return false;
int solve(){
  int ans = 0;
  for(int i=1;i<=n;i++)</pre>
  if(not lnk[i]){
    stp++; ans += dfs(i);
   }
  return ans;
}
} graph;
      Minimum Weight Matching (Clique version)
```

```
struct Graph {
 // 0-base (Perfect Match)
int n, edge[MXN][MXN];
```

```
int match[MXN], dis[MXN], onstk[MXN];
  vector<int> stk;
  void init(int _n) {
   n = _n;
   for (int i=0; i<n; i++)</pre>
    for (int j=0; j<n; j++)</pre>
     edge[i][j] = 0;
  void set_edge(int u, int v, int w) {
  edge[u][v] = edge[v][u] = w;
  bool SPFA(int u){
  if (onstk[u]) return true;
   stk.PB(u);
   onstk[u] = 1;
   for (int v=0; v<n; v++){
    if (u != v && match[u] != v && !onstk[v]){
     int m = match[v];
     if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
      dis[m] = dis[u] - edge[v][m] + edge[u][v];
      onstk[v] = 1;
      stk.PB(v)
      if (SPFA(m)) return true;
      stk.pop_back();
      onstk[v] = 0;
    }
  onstk[u] = 0;
   stk.pop_back();
   return false;
  int solve() {
  // find a match
   for (int i=0; i<n; i+=2){</pre>
   match[i] = i+1;
   match[i+1] = i;
   while (true){
    int found = 0;
    for (int i=0; i<n; i++)</pre>
     dis[i] = onstk[i] = 0;
    for (int i=0; i<n; i++){</pre>
     stk.clear()
     if (!onstk[i] && SPFA(i)){
      found = 1:
      while (SZ(stk)>=2){
       int u = stk.back(); stk.pop_back();
       int v = stk.back(); stk.pop_back();
       match[u] = v;
       match[v] = u;
   if (!found) break;
   int ret = 0:
   for (int i=0; i<n; i++)
    ret += edge[i][match[i]];
   return ret>>1;
} graph;
4.5 Flow Models
   · Maximum/Minimum flow with lower bound / Circulation problem
```

- 1. Construct super source ${\cal S}$ and sink ${\cal T}$.
- 2. For each edge (x,y,l,u), connect $x\to y$ with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum
- of incoming lower bounds and the sum of outgoing lower bounds. 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect
- $v \to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect t o s with capacity ∞ and let the flow from S to T be f'. If $f+f'
 eq \sum_{v\in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
- 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph(X,Y)

```
Redirect every edge: y \to x if (x,y) \in M, x \to y otherwise.
      2. DFS from unmatched vertices in X 3. x \in X is chosen iff x is unvisited.
      4. y \in Y is chosen iff y is visited.
· Minimum cost cyclic flow
      1. Consruct super source {\cal S} and sink {\cal T}
      2. For each edge (x, y, c), connect x \to y with (cost, cap) = (c, 1) if
          c>0, otherwise connect y 	o x with (cost, cap)=(-c,1)
      3. For each edge with c < 0, sum these cost as K, then increase d(y)
         by 1, decrease d(x) by 1
      4. For each vertex v with d(v) > 0, connect S \to v with (cost, cap) =
         (0, d(v))
      5. For each vertex v with d(v) < 0, connect v \to T with (cost, cap) =
      6. Flow from S to T, the answer is the cost of the flow C+K
· Maximum densitu induced subgraph
      1. Binary search on answer, suppose we're checking answer {\cal T}
      2. Construct a max flow model, let K be the sum of all weights
      3. Connect source s \to v , v \in G with capacity K
      4. For each edge (u,v,w) in G, connect u \to v and v \to u with capacity
      5. For v \in \mathit{G}, connect it with sink v \to t with capacity K + 2T -
         (\sum_{e \in E(v)} w(e)) - 2w(v)
      6. T is a valid answer if the maximum flow f < K|V|
· Minimum weight edge cover
      1. For each v \in V create a copy v', and connect u' \to v' with weight
         w(u,v).
```

Project selection problem

cheapest edge incident to v.

1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge

2. Connect v
ightarrow v' with weight $2\mu(v)$, where $\mu(v)$ is the cost of the

- (v, t) with capacity $-p_v$. 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

3. Find the minimum weight perfect matching on G'.

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with ca-
- 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4.6 Dinic

```
class Dinic{
private:
using CapT = int64_t;
struct Edge{
  int to, rev;
  CapT cap;
};
int n, st, ed;
vector<vector<Edge>> G;
vector<int> lv, idx;
bool BFS(){
 fill(lv.begin(), lv.end(), -1);
  queue<int> bfs;
  bfs.push(st);
  lv[st] = 0;
  while(!bfs.empty()){
  int u = bfs.front(); bfs.pop();
   for(auto e: G[u]){
    if(e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
    lv[e.to] = lv[u] + 1;
    bfs.push(e.to);
  return (lv[ed]!=-1);
CapT DFS(int u, CapT f){
  if(u == ed) return f;
  CapT ret = 0;
  for(int& i = idx[u]; i < (int)G[u].size(); ++i){</pre>
   auto& e = G[u][i];
   if(e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
   CapT nf = DFS(e.to, min(f, e.cap));
   ret += nf; e.cap -= nf; f -= nf;
   G[e.to][e.rev].cap += nf;
   if(f == 0) return ret;
```

```
if(ret == 0) lv[u] = -1;
  return ret;
public:
 void init(int n_, int st_, int ed_){
  n = n_, st = st_, ed = ed_;
  G.resize(n); lv.resize(n);
  fill(G.begin(), G.end(), vector<Edge>());
 void add_edge(int u, int v, CapT c){
  G[u].push_back({v, (int)G[v].size(), c});
  G[v].push_back({u, ((int)G[u].size())-1, 0});
 CapT max_flow(){
  CapT ret = 0;
  while(BFS()){
   idx.assign(n, 0);
   CapT f = DFS(st, numeric_limits<CapT>::max());
   ret += f:
   if(f == 0) break;
  return ret;
} flow;
4.7
      Minimum Cost Maximum Flow
class MiniCostMaxiFlow{
 using CapT = int;
 using WeiT = int64_t;
 using PCW = pair<CapT,WeiT>;
 static constexpr CapT INF_CAP = 1 << 30;</pre>
 static constexpr WeiT INF_WEI = 1LL<<60;</pre>
private:
 struct Edge{
  int to, back;
  WeiT wei;
  CapT cap
  Edge() {}
  Edge(int a,int b,WeiT c,CapT d):
   to(a),back(b),wei(c),cap(d)
  {}
 };
 int ori, edd;
 vector<vector<Edge>> G;
 vector<int> fa, wh;
 vector<bool> inq;
 vector<WeiT> dis;
 PCW SPFA(){
  fill(inq.begin(),inq.end(),false);
  fill(dis.begin(), dis.end(), INF_WEI);
  queue<int> qq; qq.push(ori);
  dis[ori]=0;
  while(!qq.empty()){
   int u=qq.front();qq.pop();
   inq[u] = 0;
   for(int i=0;i<SZ(G[u]);++i){</pre>
    Edge e=G[u][i];
    int v=e.to;
    WeiT d=e.wei;
    if(e.cap<=0||dis[v]<=dis[u]+d)
     continue
    dis[v]=dis[u]+d;
    fa[v]=u, wh[v]=i;
    if(inq[v]) continue;
    qq.push(v);
    inq[v]=1;
  if(dis[edd]==INF_WEI)
   return {-1,-1};
  CapT mw=INF_CAP;
  for(int i=edd;i!=ori;i=fa[i])
   mw=min(mw,G[fa[i]][wh[i]].cap);
  for (int i=edd;i!=ori;i=fa[i]){
   auto &eg=G[fa[i]][wh[i]];
   eg.cap-=mw;
   G[eg.to][eg.back].cap+=mw;
  return {mw,dis[edd]};
 }
```

public:

void init(int a,int b,int n){

```
ori=a,edd=b;
  G.clear();G.resize(n);
  fa.resize(n);wh.resize(n);
  inq.resize(n); dis.resize(n);
 void add_edge(int st,int ed,WeiT w,CapT c){
 G[st].emplace_back(ed,SZ(G[ed]),w,c);
 G[ed].emplace_back(st,SZ(G[st])-1,-w,0);
PCW solve(){
 /* might modify to
 cc += ret.first * ret.second
 or
 ww += ret.first * ret.second
 CapT cc=0; WeiT ww=0;
 while(true)
  PCW ret=SPFA();
  if(ret.first==-1) break;
   cc+=ret.first;
  ww+=ret.second;
  }
 return {cc,ww};
}
} mcmf;
4.8 Global Min-Cut
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
pair<int, int> phase(int n) {
memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
int s = -1, t = -1;
while (true) {
  int c = -1;
  for (int i = 0; i < n; ++i) {</pre>
  if (del[i] || v[i]) continue;
  if (c == -1 \mid | g[i] > g[c]) c = i;
 if (c == -1) break;
 v[s = t, t = c] = true;
 for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
   g[i] += w[c][i];
return make_pair(s, t);
int mincut(int n) {
int cut = 1e9;
memset(del, false, sizeof(del));
for (int i = 0; i < n - 1; ++i)
 int s, t; tie(s, t) = phase(n);
 del[t] = true; cut = min(cut, g[t]);
 for (int j = 0; j < n; ++j) {
 w[s][j] += w[t][j]; w[j][s] += w[j][t];</pre>
return cut;
5
    Math
```

5.1 Prime Table

```
\begin{array}{c} 1002939109, 1020288887, 1028798297, 1038684299, \\ 1041211027, 1051762951, 1058585963, 1063020809, \\ 1147930723, 1172520109, 1183835981, 1187659051, \\ 1241251303, 1247184097, 1255940849, 1272759031, \\ 1287027493, 1288511629, 1294632499, 1312650799, \\ 1868732623, 1884198443, 1884616807, 1885059541, \\ 1909942399, 1914471137, 1923951707, 1925453197, \\ 1979612177, 1980446837, 1989761941, 2007826547, \\ 2008033571, 2011186739, 2039465081, 2039728567, \\ 2093735719, 2116097521, 2123852629, 2140170259, \\ 3148478261, 3153064147, 3176351071, 3187523093, \\ 3196772239, 3201312913, 3203063977, 3204840059, \\ 3210224309, 3213032591, 3217689851, 3218469083, \\ 3219857533, 3231880427, 3235951699, 3273767923, \\ 3276188869, 3277183181, 3282463507, 3285553889, \\ 3319309027, 3327005333, 3327574903, 3341387953, \\ 3373293941, 3380077549, 3380892997, 3381118801 \end{array}
```

```
5.2 \lfloor \frac{n}{i} \rfloor Enumeration
T_0 = 1, \overline{T_{i+1}} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor
5.3 ax+by=gcd
// ax+ny = 1, ax+ny == ax == 1 \pmod{n}
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
 if (y == 0) g=x,a=1,b=0;
 else exgcd(y,x%y,g,b,a),b=(x/y)*a;
5.4 Pollard Rho
// does not work when n is prime
// return any non-trivial factor
llu pollard_rho(llu n){
 static auto f=[](llu x,llu k,llu m){
  return add(k,mul(x,x,m),m);
 if (!(n&1)) return 2;
 mt19937 rnd(120821011);
 while(true){
  llu y=2,yy=y,x=rnd()%n,t=1;
  for(llu sz=2;t==1;sz<<=1) {</pre>
   for(llu i=0;i<sz;++i){</pre>
    if(t!=1)break;
    yy=f(yy,x,n);
    t=gcd(yy>y?yy-y:y-yy,n);
   y=yy;
  if(t!=1&&t!=n) return t;
      Pi Count (Linear Sieve)
static constexpr int N = 1000000 + 5;
1ld pi[N];
vector<int> primes;
bool sieved[N]
11d cube_root(11d x){
 lld s=cbrt(x-static_cast<long double>(0.1));
 while(s*s*s <= x) ++s;
 return s-1;
1ld square_root(lld x){
 lld s=sqrt(x-static_cast<long double>(0.1));
 while(s*s <= x) ++s;
 return s-1;
void init(){
 primes.reserve(N);
 primes.push_back(1);
 for(int i=2;i<N;i++) {</pre>
  if(!sieved[i]) primes.push_back(i);
  pi[i] = !sieved[i] + pi[i-1];
for(int p: primes) if(p > 1) {
   if(p * i >= N) break;
    sieved[p * i] = true;
    if(p % i == 0) break;
11d phi(11d m, 11d n) {
 static constexpr int MM = 80000, NN = 500;
 static lld val[MM][NN];
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
 if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
 return ret;
11d pi_count(11d);
11d P2(11d m, 11d n) {
 11d sm = square_root(m), ret = 0;
 for(lld i = n+1;primes[i]<=sm;i++)</pre>
  ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
 return ret;
11d pi_count(11d m) {
 if(m < N) return pi[m];</pre>
 11d n = pi_count(cube_root(m));
```

```
return phi(m, n) + n - 1 - P2(m, n);
                                                              bool notprime[N];
                                                              11d phi[N];
                                                              void euler_sieve(int n){
5.6 Range Sieve
                                                               for(int i=2;i<n;i++){</pre>
                                                                if(!notprime[i]){
const int MAX_SQRT_B = 50000;
                                                                 primes.push_back(i); phi[i] = i-1;
const int MAX_L = 200000 + 5;
                                                                for(auto j: primes){
bool is_prime_small[MAX_SQRT_B];
                                                                 if(i*j >= n) break;
bool is_prime[MAX_L];
                                                                 notprime[i*j] = true;
                                                                 phi[i*j] = phi[i] * phi[j];
void sieve(lld 1, lld r){
                                                                 if(i \% j == 0){
                                                                  phi[i*j] = phi[i] * j;
 for(lld i=2;i*i<r;i++) is_prime_small[i] = true;</pre>
                                                                   break;
 for(lld i=1;i<r;i++) is_prime[i-1] = true;
if(l==1) is_prime[0] = false;</pre>
 for(lld i=2;i*i<r;i++){</pre>
  if(!is_prime_small[i]) continue;
                                                              }
  for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;
for(lld j=std::max(2LL, (l+i-1)/i)*i;j<r;j+=i)</pre>
                                                                      Gauss Elimination
                                                              5.10
    is_prime[j-l]=false;
                                                              void gauss(vector<vector<double>> &d) {
}
                                                                int n = d.size(), m = d[0].size();
                                                                for (int i = 0; i < m; ++i) {
5.7 Miller Rabin
                                                                  int p = -1;
                                                                   for (int j = i; j < n; ++j) {
  if (fabs(d[j][i]) < eps) continue;</pre>
bool isprime(llu x){
 static llu magic[]={2,325,9375,28178,\
                                                                     if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
           450775,9780504,1795265022};
 static auto witn=[](llu a,llu u,llu n,int t){
                                                                  if (p == -1) continue;
  a = mpow(a,u,n);
                                                                   for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
  if (!a)return 0;
                                                                   for (int j = 0; j < n; ++j) {
  while(t--){
                                                                     if (i == j) continue;
   1lu a2=mul(a,a,n);
                                                                     double z = d[j][i] / d[i][i];
   if(a2==1 && a!=1 && a!=n-1)
                                                                     for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
    return 1:
   a = a2:
                                                                }
  }
                                                              }
  return a!=1;
 };
                                                              5.11
                                                                     Fast Fourier Transform
 if(x<2)return 0;</pre>
 if(!(x&1))return x==2;
 llu x1=x-1; int t=0;
                                                                polynomial multiply:
 while(!(x1&1))x1>>=1,t++;
                                                                DFT(a, len); DFT(b, len);
 for(llu m:magic)if(witn(m,x1,x,t))return 0;
                                                                for(int i=0;i<len;i++) c[i] = a[i]*b[i];
 return 1;
                                                                iDFT(c, len);
                                                                (len must be 2^k and = 2^m(max(a, b)))
                                                                Hand written Cplx would be 2x faster
5.8 Inverse Element
                                                              Cplx omega[2][N];
// x's inverse mod k
                                                              void init_omega(int n) {
long long GetInv(long long x, long long k){
                                                               static constexpr llf PI=acos(-1);
 // k is prime: euler_(k)=k-1
                                                               const llf arg=(PI+PI)/n;
 return qPow(x, euler_phi(k)-1);
                                                               for(int i=0;i<n;++i)</pre>
// if you need [1, x] (most use: [1, k-1]
                                                                omega[0][i]={cos(arg*i),sin(arg*i)};
void solve(int x, long long k){
                                                               for(int i=0;i<n;++i)</pre>
                                                                omega[1][i]=conj(omega[0][i]);
 inv[1] = 1;
 for(int i=2;i<x;i++)</pre>
                                                              void tran(Cplx arr[],int n,Cplx omg[]) {
  inv[i] = ((long long)(k - k/i) * inv[k % i]) % k;
                                                               for(int i=0, j=0;i<n;++i){</pre>
                                                                if(i>j)swap(arr[i],arr[j]);
    Euler Phi Function
                                                                for(int l=n>>1;(j^=1)<1;l>>=1);
                                                               for (int l=2;l<=n;l<<=1){
  extended euler:
                                                                int m=1>>1;
  a^b mod p
                                                                for(auto p=arr;p!=arr+n;p+=1){
  if gcd(a, p)==1: a^(b%phi(p))
                                                                 for(int i=0;i<m;++i){</pre>
  elif b < phi(p): a^b mod p
                                                                  Cplx t=omg[n/l*i]*p[m+i];
  else a^(b%phi(p) + phi(p))
                                                                  p[m+i]=p[i]-t; p[i]+=t;
lld euler_phi(int x){
 11d r=1;
 for(int i=2;i*i<=x;++i){</pre>
  if(x%i==0){
                                                              void DFT(Cplx arr[],int n){tran(arr,n,omega[0]);}
   x/=i; r*=(i-1);
                                                              void iDFT(Cplx arr[],int n){
   while(x%i==0){
                                                               tran(arr,n,omega[1]);
    x/=i; r*=i;
                                                               for(int i=0;i<n;++i) arr[i]/=n;</pre>
  }
                                                                     High Speed Linear Recurrence
 if(x>1) r*=x-1;
                                                              #define mod 998244353
 return r;
                                                              const int N=1000010;
                                                              int n,k,m,f[N],h[N],a[N],b[N],ib[N];
vector<int> primes;
```

```
int pw(int x,int y){
                                                               n=rd();k=rd();
 int re=1;
                                                               for(int i=1;i<=k;++i)f[i]=(mod+rd())%mod;</pre>
                                                               for(int i=0;i<k;++i)h[i]=(mod+rd())%mod;</pre>
 if(y<0)y+=mod-1;
                                                               for(int i=a[k]=b[k]=1;i<=k;++i)</pre>
 while(y){
  if(y&1)re=(11)re*x%mod;
                                                                a[k-i]=b[k-i]=(mod-f[i])%mod;
  y>>=1; x=(11)x*x%mod;
                                                               int len=1;while(len<=(k<<1))len<<=1;</pre>
                                                               reverse(a,a+k+1);
 return re;
                                                               poly::inv(a,ib,len);
                                                               poly::cls(ib,k+1,len);
void inc(int&x,int y){x+=y;if(x>=mod)x-=mod;}
                                                               poly::ntt(b,len,1);
namespace poly{
                                                               poly::ntt(ib,len,1);
 const int G=3;
                                                               poly::pow(a,n);
 int rev[N],L;
                                                               int ans=0;
 void ntt(int*A,int len,int f){
                                                               for(int i=0;i<k;++i)inc(ans,(11)a[i]*h[i]%mod);</pre>
                                                               printf("%d\n",ans);
  for(L=0;(1<<L)<len;++L);
  for(int i=0;i<len;++i){</pre>
                                                               return 0;
   rev[i]=(rev[i>>1]>>1)|((i&1)<<(L-1));
   if(i<rev[i])swap(A[i],A[rev[i]]);</pre>
                                                              5.13
                                                                    Chinese Remainder
  for(int i=1;i<len;i<<=1){</pre>
                                                              lld crt(lld ans[], lld pri[], int n){
   int wn=pw(G, f*(mod-1)/(i<<1));</pre>
                                                               11d M = 1, ret = 0;
   for(int j=0;j<len;j+=i<<1){</pre>
                                                               for(int i=0;i<n;i++) M *= pri[i];
    int w=1;
                                                               for(int i=0;i<n;i++){</pre>
    1ld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
     int x=A[j+k],y=(11)w*A[j+k+i]%mod;
                                                                ret += (ans[i]*(M/pri[i])%M * iv)%M;
     A[j+k]=(x+y) \mod A[j+k+i]=(x-y+mod) \mod ;
                                                                ret %= M:
   }
                                                               return ret;
                                                              }
  if(!~f){
                                                              /*
   int iv=pw(len,mod-2);
                                                              Another:
   for(int i=0;i<len;++i)A[i]=(11)A[i]*iv%mod;</pre>
                                                              x = a1 \% m1
                                                              x = a2 \% m2
 }
                                                              g = gcd(m1, m2)
 void cls(int*A,int l,int r){
                                                              assert((a1-a2)%g==0)
  for(int i=1;i<r;++i)A[i]=0;}</pre>
                                                              [p, q] = exgcd(m2/g, m1/g)
 void cpy(int*A,int*B,int 1){
                                                              return a2+m2*(p*(a1-a2)/g)
  for(int i=0;i<1;++i)A[i]=B[i];}</pre>
                                                              0 <= x < lcm(m1, m2)
 void inv(int*A,int*B,int 1){
  if(l==1){B[0]=pw(A[0], mod-2); return;}
  static int t[N];
                                                              5.14 Berlekamp Massey
  int len=l<<1;
                                                              // x: 1-base, p[]: 0-base
  inv(A,B,l>>1);
                                                              template<size_t N>
  cpy(t,A,1);cls(t,1,len);
                                                              vector<llf> BM(llf x[N], size_t n){
  ntt(t,len,1);ntt(B,len,1);
                                                                size_t f[N]={0},t=0;llf d[N];
  for(int i=0;i<len;++i)</pre>
                                                                vector<llf> p[N];
   B[i]=(11)B[i]*(2-(11)t[i]*B[i]%mod+mod)%mod;
                                                                for(size_t i=1,b=0;i<=n;++i) {</pre>
  ntt(B,len,-1);cls(B,l,len);
                                                                  for(size_t j=0;j<p[t].size();++j)</pre>
                                                                    d[i]+=x[i-j-1]*p[t][j];
 void pmod(int*A){
                                                                  if(abs(d[i]-=x[i])<=EPS)continue;</pre>
  static int t[N];
                                                                  f[t]=i;if(!t){p[++t].resize(i);continue;}
  int l=k+1,len=1;while(len<=(k<<1))len<<=1;</pre>
                                                                  vector<llf> cur(i-f[b]-1);
  cpy(t, A, (k<<1)+1);
                                                                  11f k=-d[i]/d[f[b]];cur.PB(-k);
  reverse(t,t+(k<<1)+1);
                                                                  for(size_t j=0;j<p[b].size();j++)</pre>
  cls(t,1,len);
                                                                    cur.PB(p[b][j]*k);
  ntt(t,len,1);
                                                                  if(cur.size()<p[t].size())cur.resize(p[t].size());</pre>
  for(int i=0;i<len;++i)t[i]=(11)t[i]*ib[i]%mod;</pre>
                                                                  for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];</pre>
  ntt(t,len,-1);
                                                                  if(i-f[b]+p[b].size()>=p[t].size()) b=t;
  cls(t,1,len)
                                                                  p[++t]=cur;
  reverse(t,t+1);
  ntt(t,len,1)
                                                                return p[t];
  for(int i=0;i<len;++i)t[i]=(11)t[i]*b[i]%mod;</pre>
                                                              }
  ntt(t,len,-1);
  cls(t,1,len);
                                                              5.15 NTT
  for(int i=0;i<k;++i)A[i]=(A[i]-t[i]+mod)%mod;</pre>
  cls(A,k,len);
                                                              // Remember coefficient are mod P
                                                              /* p=a*2^n+1
 void pow(int*A,int n){
                                                                    2<sup>n</sup>
                                                                n
                                                                                                  root
  if(n==1) {cls(A,0,k+1);A[1]=1;return;}
                                                                16 65536
                                                                                 65537
                                                                                             1
                                                                                                  3
                                                                                                  3 */
                                                                                 7340033
  pow(A, n>>1);
                                                                20 1048576
                                                                                             7
  int len=1; while(len<=(k<<1))len<<=1;</pre>
                                                              // (must be 2<sup>k</sup>)
                                                              template<LL P, LL root, int MAXN>
  ntt(A,len,1);
  for(int i=0;i<len;++i)A[i]=(11)A[i]*A[i]%mod;</pre>
                                                              struct NTT{
                                                               static LL bigmod(LL a, LL b) {
  ntt(A,len,-1);
  pmod(A)
                                                                LL res = 1;
  if(n&1){
                                                                for (LL bs = a; b; b >>= 1, bs = (bs * bs) % P)
   for(int i=k;i;--i)A[i]=A[i-1];A[0]=0;
                                                                 if(b&1) res=(res*bs)%P;
   pmod(A);
                                                                return res;
                                                               static LL inv(LL a, LL b) {
                                                                if(a==1)return 1;
int main(){
                                                                return (((LL)(a-inv(b%a,a))*b+1)/a)%b;
```

```
return res;
LL omega[MAXN+1];
                                                                   Poly Modulo(const Poly &a, const Poly &b) {
NTT() +
  omega[0] = 1;
                                                                      if (a.size() < b.size()) return a;</pre>
  LL r = bigmod(root, (P-1)/MAXN);
                                                                      auto dv = Multiply(Divide(a, b), b);
  for (int i=1; i<=MAXN; i++)</pre>
                                                                      assert(dv.size() == a.size());
                                                                      for (int i = 0; i < dv.size(); ++i)
  dv[i] = (a[i] + kMod - dv[i]) % kMod;</pre>
   omega[i] = (omega[i-1]*r)%P;
 // n must be 2<sup>k</sup>
                                                                      while (!dv.empty() && dv.back() == 0) dv.pop_back();
void tran(int n, LL a[], bool inv_ntt=false){
  int basic = MAXN / n , theta = basic;
                                                                      return dv;
  for (int m = n; m >= 2; m >>= 1) {
                                                                   Poly Integral(const Poly &f) {
   int mh = m >> 1;
                                                                      int n = f.size();
   for (int i = 0; i < mh; i++) {
                                                                      VI res(n + 1);
    LL w = omega[i*theta%MAXN];
                                                                      for (int i = 0; i < n; ++i)
    for (int j = i; j < n; j += m) {</pre>
                                                                        res[i+1] = 1LL * f[i] * fpow(i + 1, kMod - 2)%kMod;
     int k = j + mh;
                                                                      return res;
     LL x = a[j] - a[k];
     if (x < 0) x += P;
                                                                   Poly Evaluate(const Poly &f, const VI &x) {
     a[j] += a[k];
                                                                      if (x.empty()) return Poly();
     if (a[j] > P) a[j] -= P;
                                                                      int n = x.size();
     a[k] = (w * x) % P;
                                                                      vector<Poly> up(n * 2);
                                                                      for (int i = 0; i < n; ++i) up[i+n] = {kMod-x[i], 1};
for (int i = n - 1; i > 0; --i)
                                                                      up[i] = Multiply(up[i * 2], up[i * 2 + 1]);
   theta = (theta * 2) % MAXN;
                                                                      vector<Poly> down(n * 2)
  int i = 0:
                                                                      down[1] = Modulo(f, up[1]);
  for (int j = 1; j < n - 1; j++) {
                                                                      for (int i = 2; i < n * 2; ++i)</pre>
   for (int k = n >> 1; k > (i ^= k); k >>= 1); if (j < i) swap(a[i], a[j]);
                                                                      down[i] = Modulo(down[i >> 1], up[i]);
                                                                      VI y(n);
                                                                      for (int i = 0; i < n; ++i) y[i] = down[i + n][0];
  if (inv_ntt) {
                                                                      return y;
   LL ni = inv(n,P);
   reverse( a+1 , a+n );
                                                                   Poly Interpolate(const VI &x, const VI &y) {
   for (i = 0; i < n; i++)
                                                                      int n = x.size();
    a[i] = (a[i] * ni) % P;
                                                                      vector<Poly> up(n * 2);
                                                                      for (int i = 0; i < n; ++i) up[i+n] = {kMod-x[i], 1};
                                                                      for (int i = n - 1; i > 0; --i)
up[i] = Multiply(up[i * 2], up[i * 2 + 1]);
const LL P=2013265921, root=31;
                                                                      VI a = Evaluate(Derivative(up[1]), x);
                                                                      for (int i = 0; i < n; ++i)
a[i] = 1LL * y[i] * fpow(a[i], kMod - 2) % kMod;
const int MAXN=4194304;
NTT<P, root, MAXN> ntt;
                                                                      vector<Poly> down(n * 2);
5.16 Polynomial Operations
                                                                      for (int i = 0; i < n; ++i) down[i + n] = {a[i]};
                                                                      for (int i = n - 1; i > 0; --i) {
  auto lhs = Multiply(down[i * 2], up[i * 2 + 1]);
using VI = vector<int>;
Poly Inverse(Poly f) {
                                                                        auto rhs = Multiply(down[i * 2 + 1], up[i * 2]);
  int n = f.size()
                                                                        assert(lhs.size() == rhs.size());
  Poly q(1, fpow(f[0], kMod - 2));
  for (int s = 2;; s <<= 1) {
   if (f.size() < s) f.resize(s);</pre>
                                                                        down[i].resize(lhs.size());
                                                                        for (int j = 0; j < lhs.size(); ++j)
  down[i][j] = (lhs[j] + rhs[j]) % kMod;</pre>
    Poly fv(f.begin(), f.begin() + s);
    Poly fq(q.begin(), q.end());
fv.resize(s + s); fq.resize(s + s);
                                                                      return down[1];
    ntt::Transform(fv, s + s);
    ntt::Transform(fq, s + s);
for (int i = 0; i < s + s; ++i)
   fv[i] = 1LL * fv[i] * fq[i]%kMod * fq[i]%kMod;</pre>
                                                                   Poly Log(Poly f) {
                                                                      int n = f.size();
                                                                      if (n == 1) return {0};
                                                                      auto d = Derivative(f);
    ntt::InverseTransform(fv, s + s);
    Poly res(s);
                                                                      f.resize(n - 1);
    for (int i = 0; i < s; ++i) {
                                                                      d = Multiply(d, Inverse(f));
       res[i] = kMod - fv[i];
                                                                      d.resize(n - 1)
       if (i < (s >> 1)) {
                                                                      return Integral(d);
         int v = 2 * q[i] % kMod;
         (res[i] += v) >= kMod ? res[i] -= kMod : 0;
                                                                   Poly Exp(Poly f) {
                                                                      int n = f.size()
    }
                                                                      Poly q(1, 1); f[0] += 1;
                                                                      for (int s = 1; s < n; s <<= 1) {
    q = res;
    if (s >= n) break;
                                                                           (f.size() < s + s) f.resize(s + s);
                                                                        Poly g(f.begin(), f.begin() + s + s);
                                                                        Poly h(q.begin(), q.end());
h.resize(s + s); h = Log(h);
  q.resize(n);
  return q;
                                                                        for (int i = 0; i < s + s; ++i)
Poly Divide(const Poly &a, const Poly &b) {
                                                                          g[i] = (g[i] + kMod - h[i]) % kMod;
                                                                        g = Multiply(g, q);
  int n = a.size(), m = b.size(), k = 2;
  while (k < n - m + 1) k <<= 1;</pre>
                                                                        g.resize(s + s); q = g;
  Poly ra(k), rb(k);
  for (int i = 0; i < min(n, k); ++i) ra[i] = a[n-1-i];
for (int i = 0; i < min(m, k); ++i) rb[i] = b[m-1-i];
                                                                      assert(q.size() >= n);
                                                                      q.resize(n);
  auto rbi = Inverse(rb);
                                                                      return q;
  auto res = Multiply(rbi, ra);
  res.resize(n - m + 1);
                                                                   Poly SquareRootImpl(Poly f) {
  reverse(res.begin(), res.end());
                                                                     if (f.empty()) return {0};
```

N = (N * f) % P;

```
int z = QuadraticResidue(f[0], kMod), n = f.size();
  constexpr int kInv2 = (kMod + 1) >> 1;
                                                                     return -1;
  if (z = -1) return \{-1\};
                                                                    }
  VIq(1, z);
                                                                    5.19 Quadratic residue
  for (int s = 1; s < n; s <<= 1) {
                                                                    struct Status{
    if (f.size() < s + s) f.resize(s + s);</pre>
    VI fq(q.begin(), q.end());
                                                                      11 x,y;
    fq.resize(s + s)
                                                                    11 w;
    VI f2 = Multiply(fq, fq);
                                                                    Status mult(const Status& a,const Status& b,ll mod){
    f2.resize(s + s);
    for (int i = 0; i < s + s; ++i)
                                                                      res.x=(a.x*b.x+a.y*b.y%mod*w)%mod;
       f2[i] = (f2[i] + kMod - f[i]) % kMod;
    f2 = Multiply(f2, Inverse(fq));
                                                                      res.y=(a.x*b.y+a.y*b.x)%mod;
    f2.resize(s + s);
    for (int i = 0; i < s + s; ++i)
                                                                    inline Status qpow(Status _base, 11 _pow, 11 _mod) {
       fq[i] = (fq[i]+kMod - 1LL*f2[i]*kInv2%kMod)%kMod;
                                                                      Status res = \{1, 0\};
                                                                      while(_pow>0){
                                                                         if(_pow&1) res=mult(res,_base,_mod);
  q.resize(n);
                                                                         _base=mult(_base,_base,_mod);
  return q;
                                                                         _pow>>=1;
Poly SquareRoot(Poly f) {
                                                                      return res;
  int n = f.size(), m = 0;
  while (m < n \&\& f[m] == 0) m++;
                                                                    inline 11 check(11 x,11 p){
  if (m == n) return VI(n);
  if (m & 1) return {-1}
                                                                      return qpow_mod(x,(p-1)>>1,p);
  auto s = SquareRootImpl(VI(f.begin() + m, f.end()));
                                                                    inline 11 get_root(11 n,11 p){
  if (s[0] == -1) return {-1};
  VI res(n);
                                                                      if(p==2) return 1;
                                                                      if(check(n,p)==p-1) return -1;
  for (int i = 0; i < s.size(); ++i) res[i + m/2]=s[i];</pre>
                                                                      11 a;
  return res;
                                                                      while(true){
                                                                         a=rand()%p;
5.17 FWT
                                                                         w=((a*a-n)%p+p)%p;
                                                                         if(check(w,p)==p-1) break;
/* xor convolution:
* x = (x0, x1) , y = (y0, y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                      Status res = \{a, 1\}
                                                                      res=qpow(res,(p+1)>>1,p);
                                                                      return res.x;
 * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
 *z = (1/2) *z'
                                                                    5.20 De-Bruijn
 * or convolution:
                                                                    int res[maxn], aux[maxn], sz;
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                                    void db(int t, int p, int n, int k) {
 * and convolution:
                                                                     if (t > n) {
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
                                                                      if (n % p == 0)
for (int i = 1; i <= p; ++i)
const LL MOD = 1e9+7;
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
                                                                         res[sz++] = aux[i];
 for( int d = 1 ; d < N ; d <<= 1 ) {</pre>
                                                                     } else {
  int d2 = d << 1;
                                                                      aux[t] = aux[t - p];
  for( int s = 0 ; s < N ; s += d2 )
                                                                      db(t + 1, p, n, k);
   for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
LL ta = x[ i ] , tb = x[ j ];</pre>
                                                                      for (int i = aux[t - p] + 1; i < k; ++i) {
                                                                       aux[t] = i;
    x[ i ] = ta+tb;
                                                                       db(t + 1, t, n, k);
    x[ j ] = ta-tb;
if( x[ i ] >= MOD ) x[ i ] -= MOD;
                                                                      }
                                                                     }
    if(x[j] < 0) x[j] += MOD;
                                                                    int de_bruijn(int k, int n) {
  // return cyclic string of len k^n s.t. every string
 if( inv )
                                                                     // of len n using k char appears as a substring.
  for( int i = 0 ; i < N ; i++ ) {
    x[ i ] *= inv( N, MOD );</pre>
                                                                     if (k == 1) {
                                                                      res[0] = 0;
   x[ i ] %= MOD;
                                                                      return 1;
}
                                                                     for (int i = 0; i < k * n; i++) aux[i] = 0;
5.18
      DiscreteLog
                                                                     db(1, 1, n, k);
// Baby-step Giant-step Algorithm
                                                                     return sz;
11d BSGS(11d P, 11d B, 11d N) {
   // find B^L = N mod P
 unordered_map<lld, int> R;
                                                                    5.21 Simplex Construction
 1ld sq = (lld)sqrt(P);
                                                                    Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that for all 1 \leq j \leq m,
 11d t = 1:
                                                                    \sum_{1 < i < n} A_{ji} x_i \leq b_j and x_i \geq 0 for all 1 \leq i \leq n.
 for (int i = 0; i < sq; i++) {
  if (t == N) return i;
if (!R.count(t)) R[t] = i;
                                                                       1. In case of minimization, let c_i^\prime = -c_i
                                                                      2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
  t = (t * B) % P;
                                                                      3. \sum_{1 < i < n} A_{ji} x_i = b_j
 11d f = inverse(t, P);
 for(int i=0;i<=sq+1;i++) {</pre>
                                                                             • \sum_{1 \le i \le n} A_{ji} x_i \le b_j
  if (R.count(N))
                                                                             • \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j
   return i * sq + R[N];
```

4. If x_i has no lower bound, replace x_i with $x_i - x_i'$

5.22 Simplex

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
double inv = 1.0 / d[r][s];
for (int i = 0; i < m + 2; ++i)
  for (int j = 0; j < n + 2; ++j)
  if (i != r && j != s)
   d[i][j] -= d[r][j] * d[i][s] * inv;
for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
int x = m + z;
while (true) {
 int s = -1;
  for (int i = 0; i <= n; ++i) {
  if (!z && q[i] == -1) continue;
  if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
  if (d[x][s] > -eps) return true;
  int r = -1
  for (int i = 0; i < m; ++i) {
  if (d[i][s] < eps) continue;
if (r == -1 || \</pre>
   d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
 pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
m = b.size(), n = c.size();
d = VVD(m + 2, VD(n + 2));
for (int i = 0; i < m; ++i)</pre>
 for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)
 p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];
q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
for (int i = 1; i < m; ++i)
  if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
 if (d[r][n + 1] < -eps) {</pre>
 pivot(r, n)
  if (!phase(1) \mid | d[m + 1][n + 1] < -eps)
  return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
  int s = min_element(d[i].begin(), d[i].end() - 1)
       - d[i].begin();
   pivot(i, s);
  }
if (!phase(0)) return VD(n, inf);
VD x(n);
for (int i = 0; i < m; ++i)
 if (p[i] < n) \times [p[i]] = d[i][n + 1];
return x;
}}
```

6 Geometry

6.1 Circle Class

```
template<typename T>
struct Circle{
  static constexpr llf EPS = 1e-8;
  Point<T> o; T r;
  vector<Point<llf>> operator&(const Circle& aa)const{
    llf d=o.dis(aa.o);
    if(d>r+aa.r+EPS || d<fabs(r-aa.r)-EPS) return {};</pre>
```

```
17
  11f dt = (r*r - aa.r*aa.r)/d, d1 = (d+dt)/2;
  Point<llf> dir = (aa.o-o); dir /= d;
  Point<llf> pcrs = dir*d1 + o;
  dt=sqrt(max(0.0L, r*r - d1*d1)), dir=dir.rot90();
  return {pcrs + dir*dt, pcrs - dir*dt};
};
6.2 Segment Class
const long double EPS = 1e-8;
template<typename T>
struct Segment{
 // p1.x < p2.x
 Line<T> base;
 Point<T> p1, p2;
 Segment(): base(Line<T>()), p1(Point<T>()), p2(Point<T</pre>
    >()){
  assert(on_line(p1, base) and on_line(p2, base));
 Segment(Line<T> _, Point<T> __, Point<T> __): base(_)
   , p1(__), p2(___){
  assert(on_line(p1, base) and on_line(p2, base));
 template<typename T2>
  Segment(const Segment<T2>& _): base(_.base), p1(_.p1)
     , p2(_.p2) {}
 typedef Point<long double> Pt;
 friend bool on_segment(const Point<T>& p, const
    Segment& 1){
  if(on_line(p, l.base))
   return (1.p1.x-p.x)*(p.x-1.p2.x)>=0 and (1.p1.y-p.y)
    *(p.y-1.p2.y)>=0;
  return false;
 friend bool have_inter(const Segment& a, const Segment
    & b){
  if(is_parallel(a.base, b.base)){
   return on_segment(a.p1, b) or on_segment(a.p2, b) or
     on_segment(b.p1, a) or on_segment(b.p2, a);
  Pt inter = get_inter(a.base, b.base);
  return on_segment(inter, a) and on_segment(inter, b);
 friend inline Pt get_inter(const Segment& a, const
    Segment& b){
  if(!have_inter(a, b)){
   return NOT_EXIST
  }else if(is_parallel(a.base, b.base)){
   if(a.p1 == b.p1){
    if(on_segment(a.p2, b) or on_segment(b.p2, a))
    return INF_P;
    else return a.p1;
   }else if(a.p1 == b.p2){
    if(on_segment(a.p2, b) or on_segment(b.p1, a))
    return INF_P;
    else return a.p1;
   }else if(a.p2 == b.p1){
    if(on_segment(a.p1, b) or on_segment(b.p2, a))
    return INF_P;
    else return a.p2;
   }else if(a.p2 == b.p2){
    if(on_segment(a.p1, b) or on_segment(b.p1, a))
    return INF_P;
    else return a.p2;
   return INF_P;
  }
  return get_inter(a.base, b.base);
 friend ostream& operator<<(ostream& ss, const Segment&
  ss<<o.base<<", "<<o.p1<<" ~ "<<o.p2;
  return ss;
template<typename T>
inline Segment<T> get_segment(const Point<T>& a, const
    Point<T>& b){
 return Segment<T>(get_line(a, b), a, b);
```

6.3 Line Class

```
const Point<long double> INF_P(-1e20, 1e20);
const Point<long double> NOT_EXIST(1e20, 1e-20);
template<typename T>
struct Line{
static constexpr long double EPS = 1e-8;
// ax+by+c = 0
T a, b, c;
Line(T _=0, T
                __=1, T ___=0): a(_), b(__), c(___){
 assert(fabs(a)>EPS or fabs(b)>EPS);}
template<typename T2>
 Line(const Line<T2>& x): a(x.a), b(x.b), c(x.c){}
typedef Point<long double> Pt;
bool equal(const Line& o, true_type) const {
 return fabs(a-o.a)<EPS &&
 fabs(b-o.b) < EPS && fabs(c-o.b) < EPS;}
bool equal(const Line& o, false_type) const {
 return a==o.a and b==o.b and c==o.c;}
bool operator==(const Line& o) const {
 return equal(o, is_floating_point<T>());}
bool operator!=(const Line& o) const {
 return !(*this == o);
friend inline bool on_line__(const Point<T>& p, const
    Line& 1, true_type){
  return fabs(1.a*p.x + 1.b*p.y + 1.c) < EPS;
friend inline bool on_line__(const Point<T>& p, const
    Line& 1, false_type){
  return 1.a*p.x + 1.b*p.y + 1.c == 0;
friend inline bool on_line(const Point<T>&p, const
    Line& 1){
 return on_line__(p, 1, is_floating_point<T>());
friend inline bool is_parallel__(const Line& x, const
    Line& y, true_type){
  return fabs(x.a*y.b - x.b*y.a) < EPS;
friend inline bool is_parallel__(const Line& x, const
    Line& y, false_type){
  return x.a*y.b == x.b*y.a;
friend inline bool is_parallel(const Line& x, const
    Line& y){
  return is_parallel__(x, y, is_floating_point<T>());
friend inline Pt get_inter(const Line& x, const Line&
   y){
 typedef long double llf;
 if(x==y) return INF_P;
  if(is_parallel(x, y)) return NOT_EXIST;
 llf delta = x.a*y.b - x.b*y.a;
  llf delta_x = x.b*y.c - x.c*y.b;
 llf delta_y = x.c*y.a - x.a*y.c;
  return Pt(delta_x / delta, delta_y / delta);
friend ostream&operator<<(ostream&ss, const Line&o){</pre>
 ss<<o.a<<"x+"<<o.b<<"y+"<<o.c<<"=0";
 return ss;
template<typename T>
inline Line<T> get_line(const Point<T>& a, const Point<</pre>
    T>& b){}
return Line<T>(a.y-b.y, b.x-a.x, (b.y-a.y)*a.x-(b.x-a.
    x)*a.y);
```

6.4 Triangle Circumcentre

```
6.5 2D Convex Hull
template<typename T>
class ConvexHull_2D{
private:
 typedef Point<T> PT;
 vector<PT> d;
 struct myhash{
  uint64_t operator()(const PT& a) const {
   uint64_t xx=0, yy=0;
   memcpy(&xx, &a.x, sizeof(a.x));
   memcpy(&yy, &a.y, sizeof(a.y));
   uint64_t ret = xx*17+yy*31;
   ret = (ret ^ (ret >> 16))*0x9E3779B1;
ret = (ret ^ (ret >> 13))*0xC2B2AE35;
   ret = ret ^ xx;
   return (ret ^ (ret << 3)) * yy;</pre>
 };
 unordered_set<PT, myhash> in_hull;
public:
 void init(){in_hull.clear();d.clear();}
 void insert(const PT& x){d.PB(x);}
 void solve(){
  sort(ALL(d), [](const PT& a, const PT& b){
   return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
  vector<PT> s(SZ(d)<<1); int o = 0;
  for(auto p: d) {
   while(o \ge 2 \& cross(p-s[o-2], s[o-1]-s[o-2]) <= 0)
    0--
   s[o++] = p;
  for(int i=SZ(d)-2, t = o+1;i>=0;i--){
   while(o = t\&cross(d[i] - s[o-2], s[o-1] - s[o-2]) <= 0)
    0--:
   s[o++] = d[i];
  }
  s.resize(o-1); swap(s, d);
  for(auto i: s) in_hull.insert(i);
 vector<PT> get(){return d;}
 bool in_it(const PT& x){
  return in_hull.find(x)!=in_hull.end();}
6.6 2D Farthest Pair
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {</pre>
 while(abs(cross(stk[i+1]-stk[i],
   stk[(pos+1)%n]-stk[i]))
   abs(cross(stk[i+1]-stk[i],
   stk[pos]-stk[i]))) pos = (pos+1)%n;
 ans = max({ans, dis(stk[i], stk[pos]),
  dis(stk[i+1], stk[pos])});
```

6.7 2D Closest Pair

```
struct cmp_y {
 bool operator()(const P& p, const P& q) const {
  return p.y < q.y;</pre>
 }
};
multiset<P, cmp_y> s;
void solve(P a[], int n) {
 sort(a, a + n, [](const P& p, const P& q) {
  return tie(p.x, p.y) < tie(q.x, q.y);</pre>
 11f d = INF; int pt = 0;
 for (int i = 0; i < n; ++i) {
  while (pt < i \text{ and } a[i].x - a[pt].x >= d)
   s.erase(s.find(a[pt++]));
  auto it = s.lower_bound(P(a[i].x, a[i].y - d));
  while (it != s.end() and it->y - a[i].y < d)
   d = min(d, dis(*(it++), a[i]));
  s.insert(a[i]);
}
```

// cross(pt, line.ed-line.st)>=0 <-> pt in half plane

```
6.8 kD Closest Pair (3D ver.)
                                                               vector< Line > lns;
                                                               deque< Line > que;
11f solve(vector<P> v) {
                                                               deque< Point > pt;
 shuffle(v.begin(), v.end(), mt19937());
                                                               double HPI() {
 // maybe could replace vector<P> with only P
                                                                sort( lns.begin(), lns.end() );
 unordered_map<lld, unordered_map<lld,
                                                                que.clear(); pt.clear();
  unordered_map<lld, vector<P>>>> m;
                                                                que.push_back( lns[ 0 ] );
 llf d = dis(v[0], v[1]);
                                                                for ( int i = 1 ; i < (int)lns.size() ; i ++ ) {</pre>
 auto Idx = [\&d] (11d x) -> 11d {
                                                                 if(!dcmp(lns[i].ang - lns[i-1].ang)) continue;
  return round(x * 2 / d) + 0.1;
                                                                 while ( pt.size() > 0 &&
 auto rebuild_m = [&m, &v, &Idx](int k) {
                                                                  dcmp(cross(lns[i].st,lns[i].ed,pt.back()))<0){</pre>
  m.clear();
                                                                  pt.pop_back();que.pop_back();
  for (int i = 0; i < k; ++i)
   m[Idx(v[i].x)][Idx(v[i].y)]
                                                                 while ( pt.size() > 0 &&
    [Idx(v[i].z)].push_back(v[i]);
                                                                  dcmp(cross(lns[i].st,lns[i].ed,pt.front()))<0){</pre>
 };
                                                                  pt.pop_front(); que.pop_front();
 rebuild_m(2);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
                                                                 pt.push_back(get_point( que.back(), lns[ i ] ));
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
                                                                 que.push_back( lns[ i ] );
     kz = Idx(v[i].z); bool found = false;
  for (int x = -2; x <= 2; ++x) {
                                                                while ( pt.size() > 0 &&
   const 11d nx = x + kx;
                                                                 dcmp(cross(que[0].st, que[0].ed, pt.back()))<0){</pre>
   if (m.find(nx) == m.end()) continue;
                                                                 que.pop_back();
   auto& mm = m[nx];
                                                                 pt.pop_back();
   for (int y = -2; y \le 2; ++y) {
    const lld ny = y + ky;
if (mm.find(ny) == mm.end()) continue;
                                                                while ( pt.size() > 0 &&
                                                                 dcmp(cross(que.back().st,que.back().ed,pt[0]))<0){</pre>
    auto& mmm = mm[ny];
                                                                 que.pop_front();
    for (int z = -2; z <= 2; ++z) {
                                                                 pt.pop_front();
     const 11d nz = z + kz;
     if (mmm.find(nz) == mmm.end()) continue;
                                                                pt.push_back(get_point(que.front(), que.back()));
     for (auto p: mmm[nz]) {
  if (dis(p, v[i]) < d) {</pre>
                                                                vector< Point > conv;
                                                                for ( int i = 0 ; i < (int)pt.size() ; i ++ )</pre>
       d = dis(p, v[i]);
                                                                 conv.push_back( pt[ i ] );
       found = true;
                                                                double ret = 0;
                                                                for ( int i = 1 ; i + 1 < (int)conv.size() ; i ++ )
                                                                 ret += abs(cross(conv[0], conv[i], conv[i + 1]));
                                                                return ret / 2.0;
  if (found) rebuild_m(i + 1);
                                                               6.11 Ternary Search on Integer
  else m[kx][ky][kz].push_back(v[i]);
                                                               int TernarySearch(int 1, int r) {
 return d;
                                                                // max value @ (1, r]
                                                                while (r - 1 > 1){
                                                                 int m = (1 + r) >> 1;
6.9 Simulated Annealing
                                                                 if (f(m) > f(m + 1)) r = m;
                                                                 else 1 = m;
11f anneal() {
 mt19937 rnd_engine( seed );
                                                                return 1+1;
 uniform_real_distribution< llf > rnd( 0, 1 );
 const 11f dT = 0.001;
 // Argument p
                                                               6.12 Minimum Covering Circle
 11f S_cur = calc( p ), S_best = S_cur;
for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
                                                               template<typename T>
 // Modify p to p_prime
const llf S_prime = calc( p_prime );
                                                               Circle<llf> MinCircleCover(const vector<PT>& pts){
                                                                 random_shuffle(ALL(pts));
                                                                 Circle<llf> c = \{pts[0], 0\};
  const llf delta_c = S_prime - S_cur
  llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
                                                                 for(int i=0;i<SZ(pts);i++){</pre>
                                                                   if(pts[i].in(c)) continue;
 if ( rnd( rnd_engine ) <= prob )</pre>
   S_{cur} = S_{prime}, p = p_{prime};
                                                                    c = {pts[i], 0};
  if ( S_prime < S_best ) // find min</pre>
                                                                   for(int j=0;j<i;j++){</pre>
   S_best = S_prime, p_best = p_prime;
                                                                      if(pts[j].in(c)) continue;
                                                                      c.o = (pts[i] + pts[j]) / 2;
                                                                      c.r = pts[i].dis(c.o)
 return S_best;
                                                                      for(int k=0;k<j;k++){</pre>
                                                                        if(pts[k].in(c)) continue;
6.10 Half Plane Intersection
                                                                        c = get_circum(pts[i], pts[j], pts[k]);
inline int dcmp ( double x ) {
                                                                   }
if( fabs( x ) < eps ) return 0;
return x > 0 ? 1 : -1;
                                                                 }
                                                                 return c;
struct Line {
Point st, ed;
                                                                      KDTree (Nearest Point)
 double ang;
 Line(Point _s=Point(), Point _e=Point()):
                                                               const int MXN = 100005;
  st(_s),ed(_e),ang(atan2(_e.y-_s.y,_e.x-_s.x)){}
                                                               struct KDTree {
inline bool operator< ( const Line& rhs ) const {
  if(dcmp(ang - rhs.ang) != 0) return ang < rhs.ang;</pre>
                                                                struct Node {
                                                                 int x,y,x1,y1,x2,y2;
                                                                 int id,f;
***R;
  return dcmp( cross( st, ed, rhs.st ) ) < 0;</pre>
                                                                } tree[MXN], *root;
```

int n;

sz = x.size();prefix[0]=0;power[0]=1;

```
LL dis2(int x1, int y1, int x2, int y2) {
                                                                  for(int i=1;i<=sz;i++)</pre>
 LL dx = x1-x2, dy = y1-y2;
                                                                   prefix[i]=add(mul(prefix[i-1], p), x[i-1]);
  return dx*dx+dy*dy;
                                                                  for(int i=1;i<=sz;i++)power[i]=mul(power[i-1], p);</pre>
 static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
                                                                 int query(int 1, int r){
 static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
                                                                  // 1-base (1, r]
                                                                  return sub(prefix[r], mul(prefix[1], power[r-1]));
 void init(vector<pair<int,int>> ip) {
  n = ip.size();
  for (int i=0; i<n; i++) {</pre>
   tree[i].id = i;
tree[i].x = ip[i].first;
                                                                7.2 Suffix Array
   tree[i].y = ip[i].second;
                                                                namespace sfxarray {
                                                                bool t[maxn * 2];
  root = build_tree(0, n-1, 0);
                                                                int hi[maxn], rev[maxn];
                                                                int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
 Node* build_tree(int L, int R, int d) {
                                                                int x[maxn], p[maxn], q[maxn * 2];
  if (L>R) return nullptr
                                                                // sa[i]: sa[i]-th suffix is the \
  int M = (L+R)/2; tree[M].f = d%2;
                                                                // i-th lexigraphically smallest suffix.
  nth_element(tree+L,tree+M,tree+R+1,d%2?cmpy:cmpx);
                                                                // hi[i]: longest common prefix \
  tree[M].x1 = tree[M].x2 = tree[M].x;
                                                                // of suffix sa[i] and suffix sa[i - 1].
  tree[M].y1 = tree[M].y2 = tree[M].y;
                                                                void pre(int *sa, int *c, int n, int z) {
  tree[M].L = build_tree(L, M-1, d+1);
                                                                 memset(sa, 0, sizeof(int) * n);
                                                                 memcpy(x, c, sizeof(int) * z);
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
   tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
                                                                void induce(int *sa,int *c,int *s,bool *t,int n,int z){
   tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
                                                                 memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
                                                                  if (sa[i] && !t[sa[i] - 1])
  tree[M].R = build_tree(M+1, R, d+1);
                                                                   sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  if (tree[M].R) {
                                                                 memcpy(x, c, sizeof(int) * z);
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
                                                                 for (int i = n - 1; i >= 0; --i)
   tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
                                                                  if (sa[i] && t[sa[i] - 1])
                                                                   sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
   tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
  }
                                                                void sais(int *s, int *sa, int *p, int *q,
                                                                 bool *t, int *c, int n, int z) {
bool uniq = t[n - 1] = true;
  return tree+M;
 int touch(Node* r, int x, int y, LL d2){
                                                                 int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
  LL dis = sqrt(d2)+1;
                                                                 memset(c, 0, sizeof(int) * z);
  if (x<r->x1-dis || x>r->x2+dis ||
                                                                 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
    y<r->y1-dis || y>r->y2+dis)
                                                                 for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
   return 0;
                                                                 if (uniq) {
  return 1;
                                                                  for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
                                                                  return;
 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
  if (!r || !touch(r, x, y, md2)) return;
                                                                 for (int i = n - 2; i >= 0; --i)
  LL d2 = dis2(r->x, r->y, x, y);
                                                                  t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  if (d2 < md2 \mid | (d2 == md2 && mID < r->id)) {
                                                                 pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i)
   mID = r -> id;
   md2 = d2;
                                                                  if (t[i] && !t[i - 1])
  }
                                                                   sa[--x[s[i]]] = p[q[i] = nn++] = i;
                                                                 induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)
  // search order depends on split dim
  if ((r->f == 0 && x < r->x) ||
    (r->f == 1 && y < r->y))
                                                                  if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
   nearest(r->L, x, y, mID, md2);
                                                                  bool neq = last < 0 || \
                                                                   memcmp(s + sa[i], s + last)
   nearest(r->R, x, y, mID, md2);
                                                                    (p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
  } else {
   nearest(r->R, x, y, mID, md2);
                                                                  ns[q[last = sa[i]]] = nmxz += neq;
   nearest(r->L, x, y, mID, md2);
                                                                 sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
                                                                 pre(sa, c, n, z);
for (int i = nn - 1; i >= 0; --i)
 int query(int x, int y) {
                                                                  sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  int id = 1029384756;
  LL d2 = 102938475612345678LL;
                                                                 induce(sa, c, s, t, n, z);
  nearest(root, x, y, id, d2);
  return id;
                                                                void build(const string &s) {
 }
                                                                 for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];
} tree;
                                                                 _s[(int)s.size()] = 0; // s shouldn't contain 0
                                                                 sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];
     Stringology
                                                                 for(int i = 0;
                                                                                 i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
7.1 Hash
                                                                 int ind = 0; hi[0] = 0;
class Hash{
                                                                 for (int i = 0; i < (int)s.size(); ++i) {</pre>
                                                                  if (!rev[i]) {
private:
 const int p = 127, q = 1051762951;
                                                                   ind = 0;
 int sz, prefix[N], power[N];
                                                                   continue;
int add(int x, int y){return x+y>=q?x+y-q:x+y;}
int sub(int x, int y){return x-y<0?x-y+q:x-y;}</pre>
                                                                  while (i + ind < (int)s.size() && \</pre>
 int mul(int x, int y){return 1LL*x*y%q;}
                                                                   s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
                                                                  hi[rev[i]] = ind ? ind-- : 0;
public:
 void init(const string &x){
```

}}

7.3 Aho-Corasick Algorithm

```
class AhoCorasick{
private:
  static constexpr int Z = 26;
  struct node{
   node *nxt[ Z ], *fail;
   vector< int > data;
   node(): fail( nullptr ) {
    memset( nxt, 0, sizeof( nxt ) );
    data.clear();
  } *rt;
  inline int Idx( char c ) { return c - 'a'; }
 public:
  void init() { rt = new node(); }
  void add( const string& s, int d ) {
   node* cur = rt;
   for ( auto c : s ) {
  if ( not cur->nxt[ Idx( c ) ] )
  cur->nxt[ Idx( c ) ] = new node();
    cur = cur->nxt[ Idx( c ) ];
   cur->data.push_back( d );
  void compile() {
   vector< node* > bfs;
   size_t ptr = 0;
   for ( int i = 0 ; i < Z ; ++ i ) {
    if ( not rt->nxt[ i ] ) {
     // uncomment 2 lines to make it DFA
     // rt->nxt[i] = rt;
     continue;
    rt->nxt[ i ]->fail = rt;
    bfs.push_back( rt->nxt[ i ] );
   while ( ptr < bfs.size() ) {</pre>
    node* u = bfs[ ptr ++ ];
for ( int i = 0 ; i < Z ; ++ i ) {
  if ( not u->nxt[ i ] ) {
      // u->nxt[i] = u->fail->nxt[i];
      continue;
     node* u_f = u->fail;
     while ( u_f ) {
      if ( not u_f->nxt[ i ] ) {
       u_f = u_f->fail; continue;
      u->nxt[ i ]->fail = u_f->nxt[ i ];
      break;
     if ( not u_f ) u->nxt[ i ]->fail = rt;
     bfs.push_back( u->nxt[ i ] );
  void match( const string& s, vector< int >& ret ) {
   node* u = rt;
   for ( auto c : s ) {
    while ( u != rt and not u->nxt[ Idx( c ) ] )
     u = u->fail;
    u = u->nxt[ Idx( c ) ];
    if ( not u ) u = rt;
    node* tmp = u;
    while ( tmp != rt ) {
     for ( auto d : tmp->data )
ret.push_back( d );
     tmp = tmp->fail;
  }
} ac;
```

7.4 Suffix Automaton

```
struct Node{
Node *green, *edge[26];
int max_len;
Node(const int _max_len)
: green(NULL), max_len(_max_len){
  memset(edge,0,sizeof(edge));
}
```

```
} *ROOT, *LAST;
void Extend(const int c) {
 Node *cursor = LAST;
 LAST = new Node((LAST->max_len) + 1);
 for(;cursor&&!cursor->edge[c]; cursor=cursor->green)
  cursor->edge[c] = LAST;
 if (!cursor)
 LAST->green = ROOT;
 else {
  Node *potential_green = cursor->edge[c];
  if((potential_green->max_len)==(cursor->max_len+1))
   LAST->green = potential_green;
  else {
//assert(potential_green->max_len>(cursor->max_len+1));
   Node *wish = new Node((cursor->max_len) + 1);
   for(;cursor && cursor->edge[c]==potential_green;
      cursor = cursor->green)
    cursor->edge[c] = wish;
   for (int i = 0; i < 26; i++)
    wish->edge[i] = potential_green->edge[i];
   wish->green = potential_green->green;
   potential_green->green = wish;
   LAST->green = wish;
 }
char S[10000001], A[10000001];
int N;
int main(){
 scanf("%d%s", &N, S);
 ROOT = LAST = new Node(0);
 for (int i = 0; S[i]; i++)
  Extend(S[i] -
                 'a');
 while (N--){
  scanf("%s", A);
  Node *cursor = ROOT;
  bool ans = true;
  for (int i = 0; A[i]; i++){
   cursor = cursor->edge[A[i] - 'a'];
   if (!cursor) {
    ans = false;
    break;
   }
 puts(ans ? "Yes" : "No");
 return 0;
}
7.5 KMP
vector<int> kmp(const string &s) {
 vector<int> f(s.size(), 0);
 /* f[i] = length of the longest prefix
   (excluding s[0:i]) such that it coincides
   with the suffix of s[0:i] of the same length */
 /* i + 1 - f[i] is the length of the
   smallest recurring period of s[0:i] */
 int k = 0;
 for (int i = 1; i < (int)s.size(); ++i) {</pre>
  while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
  if (s[i] == s[k]) ++k;
  f[i] = k;
 return f;
vector<int> search(const string &s, const string &t) {
 // return 0-indexed occurrence of t in s
 vector<int> f = kmp(t), r;
for (int i = 0, k = 0; i < (int)s.size(); ++i)</pre>
  while(k > 0 && (k==(int)t.size() \mid \mid s[i]!=t[k]))
   k = f[k - 1]
  if (s[i] == t[k]) ++k;
  if (k == (int)t.size()) r.push_back(i-t.size()+1);
 return res;
}
7.6 Z value
char s[MAXN];
```

int len,z[MAXN]

void Z_value() {
 int i,j,left,right;

```
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 z[left=right=0]=len;
 for(i=1;i<len;i++)</pre>
  j=max(min(z[i-left],right-i),0);
  for(;i+j<len&&s[i+j]==s[j];j++);
  if(i+(z[i] = j)>right) {
   right=i+z[i];
   left=i;
}
7.7
      Manacher
int z[maxn];
int manacher(const string& s) {
  string t = ".";
 for(char c:s)) t += c, t += '.';
 int 1 = 0, r = 0, ans = 0;
 for (int i = 1; i < t.length(); ++i) {
z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
  while (i - z[i] >= 0 && i + z[i] < t.length()) {
   if(t[i - z[i]] == t[i + z[i]]) ++z[i];
   else break;
  if (i + z[i] > r) r = i + z[i], l = i;
 for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);</pre>
 return ans;
7.8 Lexico Smallest Rotation
string mcp(string s){
 int n = s.length();
 s += s:
 int i=0, j=1;
 while (i<n && j<n){</pre>
  int k = 0;
  while (k < n \&\& s[i+k] == s[j+k]) k++;
  if (s[i+k] <= s[j+k]) j += k+1;
  else i += k+1;
  if (i == j) j++;
 int ans = i < n ? i : j;</pre>
 return s.substr(ans, n);
7.9 BWT
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a'
 vector<int> v[ SIGMA ];
 void BWT(char* ori, char* res){
  // make ori -> ori + ori
  // then build suffix array
 }
 void iBWT(char* ori, char* res){
 for( int i = 0 ; i < SIGMA ; i ++ )</pre>
   v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )</pre>
   v[ ori[i] - BASE ].push_back( i );
  vector<int> a;
  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )
for( auto j : v[ i ] ){</pre>
```

7.10 Palindromic Tree

a.push_back(j);

ptr = a[ptr];
}
res[len] = 0;

}

} bwt;

ori[ptr ++] = BASE + i;

res[i] = ori[a[ptr]];

```
struct palindromic_tree{
    struct node{
    int next[26],f,len;
    int cnt,num,st,ed;
    node(int l=0):f(0),len(l),cnt(0),num(0) {
        memset(next, 0, sizeof(next)); }
```

for(int i = 0 , ptr = 0 ; i < len ; i ++){</pre>

```
vector<node> st;
vector<char> s:
int last,n;
void init(){
 st.clear();s.clear();last=1; n=0;
 st.push_back(0);st.push_back(-1);
 st[0].f=1;s.push_back(-1); }
int getFail(int x){
 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
 return x;}
void add(int c){
 s.push_back(c-='a'); ++n;
 int cur=getFail(last);
 if(!st[cur].next[c]){
  int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
  st[now].num=st[st[now].f].num+1;
 last=st[cur].next[c];
 ++st[last].cnt;}
int size(){ return st.size()-2;}
} pt;
int main() {
string s; cin >> s; pt.init();
for (int i=0; i<SZ(s); i++) {</pre>
 int prvsz = pt.size(); pt.add(s[i]);
 if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [1,r]: s.substr(1, r-l+1)
 }
}
return 0;
```

8 Misc

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\dots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

 $\text{holds for all } 1 \leq k \leq n.$

8.1.5 Havel–Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W\subseteq X, |W|\leq |N_G(W)|$

8.1.7 Euler's planar graph formula

```
V - E + F = C + 1, E \le 3V - 6(?)
```

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Lucas's theorem

```
 \binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p}, \text{ where } m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0, \\ \text{and } n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0.
```

8.2 MaximumEmptyRect

```
int max_empty_rect(int n, int m, bool blocked[N][N]) {
static int mxu[2][N], me=0, he=1, ans=0;
for (int i=0;i<m;i++) mxu[he][i]=0;</pre>
for (int i=0;i<n;i++) {</pre>
 stack<PII, vector<PII>> stk;
 for (int j=0;j<m;++j) {
  if (blocked[i][j]) mxu[me][j]=0;</pre>
   else mxu[me][j]=mxu[he][j]+1;
   int la = j;
   while (!stk.empty()&&stk.top().FF>mxu[me][j]) {
    int x1 = i - stk.top().FF, x2 = i;
    int y1 = stk.top().SS, y2 = j;
    la = stk.top().SS; stk.pop();
    ans=max(ans,(x2-x1)*(y2-y1));
   if (stk.empty()||stk.top().FF<mxu[me][j])</pre>
    stk.push({mxu[me][j],la});
  while (!stk.empty()) {
   int x1 = i - stk.top().FF, x2 = i;
   int y1 = stk.top().SS-1, y2 = m-1;
   stk.pop(); ans=max(ans,(x2-x1)*(y2-y1));
 swap(me,he);
return ans;
```

8.3 DP-opt Condition

8.3.1 totally monotone (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

8.3.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

8.4 Convex 1D/1D DP

```
struct segment {
int i, 1, r
segment() {}
segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
dp[0] = 0:
deque<segment> dq; dq.push_back(segment(0, 1, n));
for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i);</pre>
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
   f(i, dq.back().1) < f(dq.back().i, dq.back().1))
    dq.pop_back();
  if (dq.size())
   int d = 1 << 20, c = dq.back().1;</pre>
   while (d >>= 1) if (c + d <= dq.back().r)
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.1 = c + 1;
  if (seg.1 <= n) dq.push_back(seg);</pre>
```

8.5 ConvexHull Optimization

```
inline lld DivCeil(lld n, lld d) { // ceil(n/d) return n / d + (((n < 0) != (d > 0)) && (n % d));
struct Line {
  static bool flag;
 lld a, b, 1, r; // y=ax+b in [1, r)
lld operator()(lld x) const { return a * x + b; }
  bool operator<(const Line& i) const {</pre>
   return flag ? tie(a, b) < tie(i.a, i.b) : 1 < i.1;</pre>
  11d operator&(const Line& i) const {
   return DivCeil(b - i.b, i.a - a);
};
bool Line::flag = true;
class ConvexHullMax {
  set<Line> L:
  public:
  ConvexHullMax() { Line::flag = true; }
  void InsertLine(lld a, lld b) { // add y = ax + b
  Line now = \{a, b, -INF, INF\};
   if (L.empty()) {
    L.insert(now);
    return;
  Line::flag = true;
   auto it = L.lower_bound(now);
   auto prv = it == L.begin() ? it : prev(it);
   if (it != L.end() && ((it != L.begin() &&
    (*it)(it->1) >= now(it->1) &&
    (*prv)(prv->r - 1) >= now(prv->r - 1)) ||
    (it == L.begin() && it->a == now.a))) return;
   if (it != L.begin())
    while (prv != L.begin() &&
     (*prv)(prv->1) <= now(prv->1))
      prv = --L.erase(prv)
    if (prv == L.begin() && now.a == prv->a)
     L.erase(prv);
   if (it != L.end())
   while (it != --L.end() &&
     (*it)(it->r) \le now(it->r))
      it = L.erase(it)
   if (it != L.begin())
   prv = prev(it);
    const_cast<Line*>(&*prv)->r=now.l=((*prv)&now);
   if (it != L.end())
   const_cast<Line*>(&*it)->l=now.r=((*it)&now);
  L.insert(it, now);
  11d Query(11d a) const { // query max at x=a
   if (L.empty()) return -INF;
  Line::flag = false;
   auto it = --L.upper_bound(\{0, 0, a, 0\});
   return (*it)(a);
};
 8.6 Josephus Problem
```

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

8.7 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
  dfn[u]=low[u]=++dfs_idx;
  for(int i=0;i<(int)init_g[u].size();i++){</pre>
```

```
int v=init_g[u][i];
                                                                  init_g[b].push_back(a);
  if(v==par[u]) continue;
  if(!dfn[v]){
                                                                 par[1]=-1;
   par[v]=u;
                                                                 tarjan(1);
                                                                 dfs(1,-1);
   tarjan(v);
   low[u]=min(low[u],low[v]);
                                                                printf("%d\n", max(dp[1][0], dp[1][1]));
   if(dfn[u]<low[v]){</pre>
                                                                 return 0;
                                                               }
    g[u].push_back(v);
    g[v].push_back(u);
                                                               8.8 DLX
                                                               struct DLX {
  }else{
   low[u]=min(low[u],dfn[v]);
                                                                  const static int maxn=210;
   if(dfn[v]<dfn[u]){</pre>
                                                                  const static int maxm=210;
    int temp_v=u;
                                                                  const static int maxnode=210*210;
    bcc_id++;
                                                                  int n, m, size, row[maxnode], col[maxnode];
                                                                  int U[maxnode], D[maxnode], L[maxnode], R[maxnode];
    while(temp_v!=v){
     g[bcc_id+n].push_back(temp_v);
                                                                  int H[maxn], S[maxm], ansd, ans[maxn];
     g[temp_v].push_back(bcc_id+n);
                                                                  void init(int _n, int _m) {
     temp_v=par[temp_v];
                                                                    n = _n, m = _m;
                                                                    for(int i = 0; i <= m; ++i) {</pre>
    g[bcc_id+n].push_back(v);
                                                                      S[i] = 0;
    g[v].push_back(bcc_id+n);
                                                                      U[i] = D[i] = i;
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
                                                                      L[i] = i-1, R[i] = i+1;
                                                                    R[L[0] = size = m] = 0;
                                                                    for(int i = 1; i <= n; ++i) H[i] = -1;
int dp[maxn][2],min_dp[2][2],tmp[2][2],tp[2];
                                                                  void Link(int r, int c) {
void dfs(int u,int fa){
                                                                    ++S[col[++size] = c];
                                                                    row[size] = r; D[size] = D[c];
U[D[c]] = size; U[size] = c; D[c] = size;
if(H[r] < 0) H[r] = L[size] = R[size] = size;</pre>
 if(u<=n){
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
   dfs(v,u);
                                                                      R[size] = R[H[r]];
   memset(tp,0x8f,sizeof tp);
                                                                      L[R[H[r]]] = size;
                                                                      L[size] = H[r];
   if(v<=n){
    tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
                                                                      R[H[r]] = size;
    tp[1]=max(
     dp[u][0]+dp[v][0]+1
                                                                  }
     dp[u][1]+max(dp[v][0],dp[v][1])
                                                                  void remove(int c) {
                                                                    L[R[c]] = L[c]; R[L[c]] = R[c];
                                                                    for(int i = D[c]; i != c; i = D[i])
  for(int j = R[i]; j != i; j = R[j]) {
    U[D[j]] = U[j];
   }else{
    tp[0]=dp[u][0]+dp[v][0];
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
                                                                        D[U[j]] = D[j];
   dp[u][0]=tp[0],dp[u][1]=tp[1];
                                                                        --S[col[j]];
 }else{
  for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                  void resume(int c) {
                                                                    L[R[c]] = c; R[L[c]] = c;
  int v=g[u][i];
                                                                    for(int i = U[c]; i != c; i = U[i])
   if(v==fa) continue;
   dfs(v,u);
                                                                      for(int j = L[i]; j != i; j = L[j]) {
                                                                        U[D[j]] = j;
  min_dp[0][0]=0;
                                                                        D[U[j]] = j
  min_dp[1][1]=1;
                                                                        ++S[col[j]];
  min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
                                                                  void dance(int d) {
   if(v==fa) continue;
                                                                    if(d>=ansd) return;
   memset(tmp,0x8f,sizeof tmp);
                                                                    if(R[0] == 0) {
   tmp[0][0]=max(
                                                                      ansd = d;
   min_dp[0][0]+max(dp[v][0],dp[v][1]),
                                                                      return;
    min_dp[0][1]+dp[v][0]
                                                                    int c = R[0];
                                                                    for(int i = R[0]; i; i = R[i])
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
   tmp[1][0]=max(
                                                                      if(S[i] < S[c]) c = i;
   \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
                                                                    remove(c);
    min_dp[1][1]+dp[v][0]
                                                                    for(int i = D[c]; i != c; i = D[i]) {
                                                                      ans[d] = row[i]
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
                                                                      for(int j = R[i]; j != i; j = R[j])
   memcpy(min_dp,tmp,sizeof tmp);
                                                                        remove(col[j]);
                                                                      dance(d+1);
  dp[u][1]=max(min_dp[0][1], min_dp[1][0]);
                                                                      for(int j = L[i]; j != i; j = L[j])
  dp[u][0]=min_dp[0][0];
                                                                        resume(col[j]);
                                                                    resume(c);
int main(){
                                                                  }
int m,a,b;
                                                               } sol;
scanf("%d%d",&n,&m);
for(int i=0;i<m;i++){
  scanf("%d%d",&a,&b);</pre>
                                                                     Tree Knapsack
                                                               int dp[N][K];PII obj[N];
 init_g[a].push_back(b);
                                                               vector<int> G[N];
```

```
void dfs(int u, int mx){
 for(int s: G[u]) {
  if(mx < obj[s].first) continue;</pre>
  for(int i=0;i<=mx-obj[s].FF;i++)</pre>
   dp[s][i] = dp[u][i];
  dfs(s, mx - obj[s].first);
  for(int i=obj[s].FF;i<=mx;i++)</pre>
   dp[u][i] = max(dp[u][i],
    dp[s][i - obj[s].FF] + obj[s].SS);
int main(){
 int n, k; cin >> n >> k;
 for(int i=1;i<=n;i++){</pre>
  int p; cin >> p;
 G[p].push_back(i);
  cin >> obj[i].FF >> obj[i].SS;
 dfs(0, k); int ans = 0;
 for(int i=0;i<=k;i++) ans = max(ans, dp[0][i]);</pre>
 cout << ans << '\n';
 return 0;
8.10
      N Queens Problem
vector< int > solve( int n ) {
 // no solution when n=2, 3
 vector< int > ret;
 if ( n % 6 == 2 ) {
 for ( int i = 2 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
  ret.push_back( 3 ); ret.push_back( 1 );
  for ( int i = 7 ; i <= n ; i += 2 )
ret.push_back( i );
  ret.push_back( 5 );
 } else if ( n % 6 == 3 ) {
  for ( int i = 4 ; i <= n ; i += 2 )
  ret.push_back( i );
  ret.push_back( 2 );
  for ( int i = 5 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
  ret.push_back( 1 ); ret.push_back( 3 );
 } else {
  for ( int i = 2 ; i <= n ; i += 2 )
   ret.push_back( i );
  for ( int i = 1 ; i <= n ; i += 2 )
   ret.push_back( i );
 return ret;
}
8.11 Aliens Optimization
long long Alien() {
  long long c = kInf;
  for (int d = 60; d >= 0; --d) {
    // cost can be negative, depending on the problem.
    if (c - (1LL << d) < 0) continue;</pre>
    long long ck = c - (1LL \ll d);
    pair<long long, int> r = check(ck);
    if (r.second == k) return r.first - ck * k;
    if (r.second < k) c = ck;
 pair<long long, int> r = check(c);
return r.first - c * k;
```