7.10Palindromic Tree . . . . . . . . . .

23

23

23

23

23

23 23 23

23

23

23

23

23

25

#### Contents 8 Misc 8.1 Theorems 8.1.1 Kirchhoff's Theorem . . . . . . . . . . . . . 1 Basic 1.1 vimrc . . . 2 Data Structure 8.1.8 Pick's theorem . . . . . . . . . . . . . . . 2 8.3.2 monge condition (concave/convex) 2.5 LiChao Segment Tree . . . . . . . . . . . . . . . . . . 2.6 Treap . . 8.5 ConvexHull Optimization . . . . . . . . . . . 3 Graph 1 Basic 1.1 vimrc se is nu rnu bs=2 ru mouse=a encoding=utf-8 se cin et ts=4 sw=4 sts=4 t\_Co=256 syn on colorscheme ron filetype indent on map <F8> <ESC>:w<CR>:!clear && g++ "%" -o "%<" -4 Matching & Flow fsanitize=address -fsanitize=undefined -g && echo 4.1 Kuhn Munkres success<CR> 9 map <F9> <ESC>:w<CR>:!clear && g++ "%" -o "%<" -02 && 10 echo success<CR> 4.4 Minimum Weight Matching (Clique version) . . . . . 10 map <F10> <ESC>:!./"%<"<CR> 4.6 Dinic . . . . 11 4.7 Minimum Cost Maximum Flow . . . . . . . . . . . . . 1.2 Pragma optimization 12 12 #pragma GCC optimize("Ofast,no-stack-protector") 12 #pragma GCC optimize("no-math-errno,unroll-loops") 12 #pragma GCC target("sse,sse2,sse3,sse3,sse4") #pragma GCC target("popent,abm,mmx,avx,tune=native") 12 12 12 13 1.3 IO Optimization 13 13 13 static inline int gc() { 13 static char buf[ 1 << 20 ], \*p = buf, \*end = buf;</pre> 13 **if** ( p == end ) { 14 end = buf + fread( buf, 1, 1 << 20, stdin ); 14 if ( end == buf ) return EOF; 14 p = buf;15 5.16Polynomial Operations . . . . . . . . . . . . . . . . 15 return \*p++; 16 16 5.19Quadratic residue . . . . . . . . . . . . . . . 16 template < typename T > 17 static inline bool gn( T &\_ ) { 17 register int c = gc(); register T \_\_ = 1; \_ = 0; while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc(); if(c == '-') { \_\_ = -1; c = gc(); } if(c == EOF) return false; 17 17 6.1 Point Class . . . . . . . . . . . . . . . . . . 17 while('0'<=c&&c<='9') \_ = \_ \* 10 + c - '0', c = gc(); 18 \_ \*= \_\_; 18 return true; 6.5 Triangle Circumcentre . . . . . . . . . . . . . . . . 19 19 template < typename T, typename ...Args > static inline bool gn( T &x, Args &...args ) 19 { return gn(x) && gn(args...); } 19 6.10Half Plane Intersection . . . . . . . 19 6.11Ternary Search on Integer . . . . . . . . . . . 20 6.12Minimum Covering Circle . . . . . . . . . . . . . . . 2 Data Structure 6.13KDTree (Nearest Point) . . . . . . . . . . . 20 2.1 Bigint 7 Stringology 20 20 21 class BigInt{ 21 private: 7.4 Suffix Automaton . . . . . . . . . . . . . . . . . 21 using lld = int\_fast64\_t; 22 #define PRINTF\_ARG PRIdFAST64 22 #define LOG\_BASE\_STR "9" 7.7 Manacher 22 static constexpr lld BASE = 1000000000; 22 7.8 Lexico Smallest Rotation . . . . . . static constexpr int LOG\_BASE = 9;

22

vector<lld> dig;

```
bool neg;
 inline int len() const { return (int) dig.size(); }
 inline int cmp_minus(const BigInt& a) const {
   if(len() == 0 && a.len() == 0) return 0;
if(neg ^ a.neg)return (int)a.neg*2 - 1;
   if(len()!=a.len())
     return neg?a.len()-len():len()-a.len();
   for(int i=len()-1;i>=0;i--) if(dig[i]!=a.dig[i])
     return neg?a.dig[i]-dig[i]:dig[i]-a.dig[i];
   return 0;
 inline void trim(){
   while(!dig.empty()&&!dig.back())dig.pop_back();
   if(dig.empty()) neg = false;
public:
 BigInt(): dig(vector<lld>()), neg(false){}
 BigInt(lld a): dig(vector<lld>())
   neg = a<0; dig.push_back(abs(a));</pre>
   trim();
 BigInt(const string& a): dig(vector<lld>()){
   assert(!a.empty()); neg = (a[0]=='-');
for(int i=((int)a.size())-1;i>=neg;i-=LOG_BASE){
     11d cur = 0;
     for(int j=min(LOG_BASE-1,i-neg);j>=0;j--)
       cur = cur*10+a[i-j]-'0';
     dig.push_back(cur);
   } trim();
 inline bool operator<(const BigInt& a)const
   {return cmp_minus(a)<0;}
 inline bool operator <= (const BigInt& a) const
   {return cmp_minus(a)<=0;}
 inline bool operator == (const BigInt& a)const
   {return cmp_minus(a)==0;}
 inline bool operator!=(const BigInt& a)const
   {return cmp_minus(a)!=0;}
 inline bool operator>(const BigInt& a)const
   {return cmp_minus(a)>0;}
 inline bool operator>=(const BigInt& a)const
   {return cmp_minus(a)>=0;}
 BigInt operator-() const {
   BigInt ret = *this;
   ret.neg ^= 1;
   return ret;
 BigInt operator+(const BigInt& a) const {
   if(neg) return -(-(*this)+(-a));
if(a.neg) return (*this)-(-a);
   int n = max(a.len(), len());
   BigInt ret; ret.dig.resize(n);
   11d pro = 0;
   for(int i=0;i<n;i++) {</pre>
     ret.dig[i] = pro;
     if(i < a.len()) ret.dig[i] += a.dig[i];</pre>
     if(i < len()) ret.dig[i] += dig[i];</pre>
     pro = 0:
     if(ret.dig[i] >= BASE) pro = ret.dig[i]/BASE;
     ret.dig[i] -= BASE*pro;
   if(pro != 0) ret.dig.push_back(pro);
   return ret;
 BigInt operator-(const BigInt& a) const {
   if(neg) return -(-(*this) - (-a));
   if(a.neg) return (*this) + (-a);
   int diff = cmp_minus(a);
   if(diff < 0) return -(a - (*this));</pre>
   if(diff == 0) return 0;
   BigInt ret; ret.dig.resize(len(), 0);
for(int i=0;i<len();i++) {</pre>
     ret.dig[i] += dig[i];
     if(i < a.len()) ret.dig[i] -= a.dig[i];</pre>
     if(ret.dig[i] < 0){
       ret.dig[i] += BASE;
       ret.dig[i+1]--;
     }
   }
   ret.trim();
   return ret;
```

```
BigInt operator*(const BigInt& a) const {
  if(!len()||!a.len()) return 0;
  BigInt ret; ret.dig.resize(len()+a.len()+1);
  ret.neg = neg ^ a.neg;
  for(int i=0;i<len();i++)</pre>
    for(int j=0;j<a.len();j++){</pre>
      ret.dig[i+j] += dig[i] * a.dig[j];
      if(ret.dig[i+j] >= BASE)
        lld x = ret.dig[i+j] / BASE;
        ret.dig[i+j+1] += x;
        ret.dig[i+j] -= x * BASE;
  ret.trim();
  return ret;
BigInt operator/(const BigInt& a) const {
  assert(a.len());
  if(len() < a.len()) return 0;</pre>
  BigInt ret; ret.dig.resize(len()-a.len()+1);
  ret.neg = a.neg;
  for(int i=len()-a.len();i>=0;i--){
    11d 1 = 0, r = BASE;
    while(r-1 > 1){
      11d \ mid = (1+r)>>1;
      ret.dig[i] = mid;
      if(ret*a<=(neg?-(*this):(*this))) 1 = mid;</pre>
      else r = mid;
    ret.dig[i] = 1;
  ret.neg ^= neg; ret.trim();
  return ret:
BigInt operator%(const BigInt& a) const {
  return (*this) - (*this) / a * a;
friend BigInt abs(BigInt a){
  a.neg = 1; return a;
friend void swap(BigInt& a, BigInt& b){
  swap(a.dig, b.dig); swap(a.neg, b.neg);
friend istream& operator>>(istream& ss, BigInt& a){
  string s; ss >> s; a = s;
  return ss;
friend ostream&operator<<(ostream&o, const BigInt&a){</pre>
  if(a.len() == 0) return o << '0';</pre>
  if(a.neg) o <<</pre>
  ss << o.dig.back();
  for(int i=a.len()-2;i>=0;i--)
    o<<setw(LOG_BASE)<<setfill('0')<<a.dig[i];
  return o:
inline void print() const {
  if(len() == 0){putchar('0');return;}
  if(neg) putchar('-');
printf("%" PRINTF_ARG, dig.back());
  for(int i=len()-2;i>=0;i--)
    printf("%0" LOG_BASE_STR PRINTF_ARG, dig[i]);
#undef PRINTF_ARG
#undef LOG_BASE_STR
```

#### 2.2 Dark Magic

```
using __gnu_pbds::splay_tree_tag;
template<typename T>
using ordered_set = __gnu_pbds::tree<T,\
    __gnu_pbds::null_type,less<T>,rb_tree_tag,\
    __gnu_pbds::tree_order_statistics_node_update>;
// find_by_order, order_of_key
template<typename A,typename B>
using hTable1=__gnu_pbds::cc_hash_table<A,B>;
template<typename A,typename B>
using hTable2=__gnu_pbds::gp_hash_table<A,B>;
```

## 2.3 Disjoint Set

```
class DJS {
private:
  vector< int > fa, sz, sv;
  vector< pair< int*, int > > opt;
  void assign( int *k, int v ) {
    opt.emplace_back( k, *k );
    *k = v;
public:
  void init( int n ) {
    fa.resize( n ); iota( fa.begin(), fa.end(), 0 );
    sz.resize( n ); fill( sz.begin(), sz.end(), 1 );
    opt.clear();
  int query(int x) {return fa[x] == x?x:query(fa[x]);}
  void merge( int a, int b ) {
    int af = query( a ), bf = query( b );
    if( af == bf ) return;
    if( sz[ af ] < sz[ bf ] ) swap( af, bf );
assign( &fa[ bf ], fa[ af ] );
assign( &sz[ af ], sz[ af ] + sz[ bf ] );</pre>
  void save() { sv.push_back( (int) opt.size() ); }
  void undo() {
    int ls = sv.back(); sv.pop_back();
    while ( ( int ) opt.size() > ls )
       pair< int*, int > cur = opt.back();
       *cur.first = cur.second;
      opt.pop_back();
  }
};
```

### 2.4 Link-Cut Tree

```
struct Node{
 Node *par, *ch[2];
  int xor_sum, v;
 bool is_rev;
 Node(int _v){
    v=xor_sum=_v;is_rev=false;
   par=ch[0]=ch[1]=nullptr;
 inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
 inline void down(){
   if(is_rev)
      if(ch[0]!=nullptr) ch[0]->set_rev();
      if(ch[1]!=nullptr) ch[1]->set_rev();
      is_rev=false;
 inline void up(){
   xor_sum=v;
    if(ch[0]!=nullptr){
      xor_sum^=ch[0]->xor_sum;
      ch[0]->par=this;
    if(ch[1]!=nullptr){
      xor_sum^=ch[1]->xor_sum;
      ch[1]->par=this;
  inline bool is_root(){
    return par==nullptr ||\
      (par->ch[0]!=this && par->ch[1]!=this);
 bool is_rch(){return !is_root() && par->ch[1]==this;}
} *node[maxn], *stk[maxn];
```

```
int top;
void to_child(Node* p,Node* c,bool dir){
  p->ch[dir]=c;
  p->up();
inline void rotate(Node* node){
  Node* par=node->par;
  Node* par_par=par->par
  bool dir=node->is_rch();
  bool par_dir=par->is_rch();
  to_child(par, node->ch[!dir], dir);
  to_child(node,par,!dir);
  if(par_par!=nullptr && par_par->ch[par_dir]==par)
    to_child(par_par, node, par_dir);
  else node->par=par_par;
inline void splay(Node* node){
  Node* tmp=node
  stk[top++]=node;
  while(!tmp->is_root()){
    tmp=tmp->par;
    stk[top++]=tmp;
  while(top) stk[--top]->down();
  for(Node *fa=node->par;
   !node->is_root();
   rotate(node), fa=node->par)
    if(!fa->is_root())
      rotate(fa->is_rch()==node->is_rch()?fa:node);
inline void access(Node* node){
  Node* last=nullptr;
  while(node!=nullptr){
    splay(node)
    to_child(node, last, true);
    last=node:
    node=node->par;
inline void change_root(Node* node){
 access(node);splay(node);node->set_rev();
inline void link(Node* x, Node* y){
  change_root(x);splay(x);x->par=y;
inline void split(Node* x, Node* y) {
  change_root(x);access(y);splay(x);
  to_child(x,nullptr,true);y->par=nullptr;
inline void change_val(Node* node,int v){
  access(node);splay(node);node->v=v;node->up();
inline int query(Node* x,Node* y){
  change_root(x);access(y);splay(y);
  return y->xor_sum;
inline Node* find_root(Node* node){
  access(node);splay(node);
  Node* last=nullptr
  while(node!=nullptr){
    node->down();last=node;node=node->ch[0];
  return last;
set<pii> dic;
inline void add_edge(int u,int v){
  if(u>v) swap(u,v)
  if(find_root(node[u])==find_root(node[v])) return;
  dic.insert(pii(u,v))
  link(node[u],node[v]);
inline void del_edge(int u,int v){
  if(u>v) swap(u,v);
  if(dic.find(pii(u,v))==dic.end()) return;
  dic.erase(pii(u,v)
  split(node[u], node[v]);
```

## 2.5 LiChao Segment Tree

|struct Line{

```
int m, k, id;
Line() : id( -1 ) {}
  Line(int a, int b, int c)
    : m(a), k(b), id(c) {}
  int at( int x ) { return m * x + k; }
class LiChao {
 private:
    int n; vector< Line > nodes;
    inline int lc( int x ) { return 2 * x + 1; }
inline int rc( int x ) { return 2 * x + 2; }
    void insert( int 1, int r, int id, Line ln ) {
      int m = (1 + r) >> 1;
      if ( nodes[ id ].id == -1 ) {
        nodes[ id ] = ln;
        return;
      bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
      if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
        atLeft ^= 1; swap( nodes[ id ], ln );
      if ( r - l == 1 ) return;
      if ( atLeft ) insert( 1, m, lc( id ), ln );
      else insert( m, r, rc( id ), ln );
    int query( int 1, int r, int id, int x ) {
      int ret = 0;
      if ( nodes[ id ].id != -1 )
        ret = nodes[ id ].at( x );
      int m = ( 1 + r ) >> 1;
      if ( r - l == 1 ) return ret;
      else if ( x < m )</pre>
        return max( ret, query( 1, m, lc( id ), x ) );
      else
        return max( ret, query( m, r, rc( id ), x ) );
    }
  public:
    void build( int n_ ) {
      n = n_; nodes.clear();
      nodes.resize( n << 2, Line() );</pre>
    void insert( Line ln ) { insert( 0, n, 0, ln ); }
    int query( int x ) { return query( 0, n, 0, x ); }
} lichao;
```

## 2.6 Treap

```
namespace Treap{
  #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
  struct node{
    int size;
    uint32_t pri;
    node *lc, *rc;
    node() : size(0), pri(rand()), lc(0), rc(0) {}
    void pull() {
      size = 1
      if ( lc ) size += lc->size;
      if ( rc ) size += rc->size;
   }
 node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
    if ( L->pri > R->pri ) {
      L->rc = merge( L->rc, R ); L->pull();
      return L;
    } else {
      R->lc = merge( L, R->lc ); R->pull();
      return R;
  void split_by_size( node*rt,int k,node*&L,node*&R ) {
    if ( not rt ) L = R = nullptr;
    else if( sz( rt->lc ) + 1 <= k ) {
      split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
      L->pull();
    } else {
      R = rt:
      split_by_size( rt->lc, k, L, R->lc );
      R->pull();
```

```
}
#undef sz
}
```

# 2.7 SparseTable

```
template < typename T, typename Cmp_ = less< T > >
class SparseTable {
private:
  vector< vector< T > > tbl;
  vector< int > lg;
  T cv(Ta, Tb) {
     return Cmp_()( a, b ) ? a : b;
public:
  void init( T arr[], int n ) {
     // 0-base
     lg.resize(n + 1);
     lg[0] = -1;
    for( int i=1 ; i<=n ; ++i ) lg[i] = lg[i>>1] + 1;
tbl.resize( lg[n] + 1 );
     tbl[ 0 ].resize( n );
     copy( arr, arr + n, tbl[ 0 ].begin() );
     for ( int i = 1 ; i <= lg[ n ] ; ++ i ) {
       int len = 1 << ( i - 1 ), sz = 1 << i;</pre>
       tbl[ i ].resize( n - sz + 1 );
for ( int j = 0 ; j <= n - sz ; ++ j )
         tbl[i][j] = cv(tbl[i-1][j], tbl[i-1][j+len]);
    }
  T query( int 1, int r ) {
     // 0-base [l, r)
     int wh = lg[r - 1], len = 1 << wh;
     return cv( tbl[ wh ][ 1 ], tbl[ wh ][ r - len ] );
};
```

#### 2.8 Linear Basis

```
struct LinearBasis {
private:
  int n, sz
  vector< llu > B;
  inline llu two( int x ){ return ( ( llu ) 1 ) << x; }</pre>
public:
  void init( int n_ ) {
    n = n_{;} B.clear(); B.resize(n); sz = 0;
  void insert( llu x ) {
    // add x into B
    for ( int i = n-1; i >= 0 ; --i ) if( two(i) & x ){
  if ( B[ i ] ) x ^= B[ i ];
      else
        B[i] = x; sz++;
         for (int j = i - 1; j >= 0;
           if( B[ j ] && ( two( j ) & B[ i ] ) )
             B[ i ] ^= B[ j ];
           if ( two( i ) & B[ j ] )
B[ j ] ^= B[ i ];
eak;
         for (int j = i + 1)
        break;
      }
    }
  }
  inline int size() { return sz; }
  bool check( llu x ) {
    // is x in span(B) ?
    for ( int i = n-1 ; i \ge 0 ; --i ) if( two(i) \& x )
      if( B[ i ] ) x ^= B[ i ];
      else return false;
    return true;
  llu kth_small(llu k) {
    /** 1-base would always > 0 **/
    /** should check it **/
    /* if we choose at least one element
       but size(B)(vectors in B)==N(original elements)
       then we can't get 0 */
    llu ret = 0;
    for ( int i = 0 ; i < n ; ++ i ) if( B[ i ] ) {
  if( k & 1 ) ret ^= B[ i ];</pre>
```

```
k >>= 1;
}
return ret;
}
base;
```

# 3 Graph

## 3.1 Euler Circuit

```
bool vis[ N ]; size_t la[ K ];
void dfs( int u, vector< int >& vec ) {
  while ( la[ u ] < G[ u ].size() ) {
    if( vis[ G[ u ][ la[ u ] ].second ] ) {
        ++ la[ u ];
        continue;
    }
    int v = G[ u ][ la[ u ] ].first;
    vis[ G[ u ][ la[ u ] ].second ] = true;
    ++ la[ u ]; dfs( v, vec );
    vec.push_back( v );
}</pre>
```

## 3.2 BCC Edge

```
class BCC{
private:
  vector< int > low, dfn;
  int cnt;
  vector< bool > bridge;
  vector< vector< PII > > G;
  void dfs( int w, int f ) {
    low[ w ] = dfn[ w ] = cnt++;
    for ( auto [ u, t ] : G[ w ] ) {
      if ( u == f ) continue;
      if ( dfn[ u ] != 0 ) {
        low[ w ] = min( low[ w ], dfn[ u ] );
      }else{
        dfs( u, w );
        low[ w ] = min( low[ w ], low[ u ] );
        if ( low[ u ] > dfn[ w ] ) bridge[ t ] = true;
    }
public:
  void init( int n, int m ) {
    G.resize(n); cnt = 0;
    fill( G.begin(), G.end(), vector< PII >() );
    bridge.clear(); bridge.resize( m );
    low.clear(); low.resize( n );
dfn.clear(); dfn.resize( n );
  void add_edge( int u, int v ) {
    // should check for multiple edge
    G[ u ].emplace_back( v, cnt );
    G[ v ].emplace_back( u, cnt ++ );
  void solve(){
    cnt = 1;
    for (int i = 0; i < n; ++i)
  if (not vis[ i ]) dfs(i, i);</pre>
  // the id will be same as insert order, 0-base
  bool is_bridge( int x ) { return bridge[ x ]; }
} bcc;
```

#### 3.3 BCC Vertex

```
class BCC{
  private:
    int n, ecnt;
    vector< vector< pair< int, int > > G;
    vector< int > low, dfn, id;
    vector< bool > vis, ap;
    void dfs( int u, int f, int d ) {
       int child = 0;
       dfn[ u ] = low[ u ] = d; vis[ u ] = true;
       for ( auto e : G[ u ] ) if ( e.first != f ) {
```

```
if ( vis[ e.first ] ) {
  low[ u ] = min( low[ u ], dfn[ e.first ] );
          } else {
            dfs( e.first, u, d + 1 ); child ++;
low[ u ] = min( low[ u ], low[ e.first ] );
            if ( low[ e.first ] >= d ) ap[ u ] = true;
       if ( u == f and child <= 1 ) ap[ u ] = false;</pre>
     void mark( int u, int idd ) {
       // really?????????
       if ( ap[ u ] ) return;
       for ( auto e : G[ u ] )
         if( id[ e.second ] != -1 ) {
            id[ e.second ] = idd;
            mark( e.first, idd );
  public:
    void init( int n_ ) {
       ecnt = 0, n = n_{-};
       G.clear(); G.resize( n );
       low.resize( n ); dfn.resize( n );
       ap.clear(); ap.resize( n );
       vis.clear(); vis.resize( n );
     void add_edge( int u, int v ) {
       G[ u ].emplace_back( v, ecnt )
       G[ v ].emplace_back( u, ecnt ++ );
     void solve() {
       for ( int i = 0 ; i < n ; ++ i )
  if ( not vis[ i ] ) dfs( i, i, 0 );</pre>
       id.resize( ecnt );
       fill( id.begin(), id.end(), -1 );
       ecnt = 0:
       for ( int i = 0 ; i < n ; ++ i )
  if ( ap[ i ] ) for ( auto e : G[ i ] )
  if( id[ e.second ] != -1 ) {</pre>
              id[ e.second ] = ecnt;
              mark( e.first, ecnt ++ );
     int get_id( int x ) { return id[ x ]; }
     int count() { return ecnt; }
     bool is_ap( int u ) { return ap[ u ]; }
} bcc:
```

# 3.4 2-SAT (SCC)

```
class TwoSat{
  private:
    int n;
    vector<vector<int>> rG,G,sccs;
    vector<int> ord,idx;
    vector<bool> vis,result;
    void dfs(int u){
      vis[u]=true
      for(int v:G[u])
        if(!vis[v]) dfs(v);
      ord.push_back(u);
    void rdfs(int u){
      vis[u]=false;idx[u]=sccs.size()-1;
      sccs.back().push_back(u);
      for(int v:rG[u])
        if(vis[v])rdfs(v);
  public:
    void init(int n_){
      n=n_;G.clear();G.resize(n);
      rG.clear();rG.resize(n)
      sccs.clear();ord.clear();
      idx.resize(n);result.resize(n);
    void add_edge(int u,int v){
      G[u].push_back(v);rG[v].push_back(u);
    void orr(int x,int y){
      if ((x^y)==1)return;
```

```
add\_edge(x^1,y);\ add\_edge(y^1,x);\\
    bool solve(){
      vis.clear();vis.resize(n);
      for(int i=0;i<n;++i)</pre>
        if(not vis[i])dfs(i);
      reverse(ord.begin(),ord.end());
      for (int u:ord){
        if(!vis[u])continue;
        sccs.push_back(vector<int>());
        rdfs(u);
      for(int i=0;i<n;i+=2)</pre>
        if(idx[i]==idx[i+1])
          return false
      vector<bool> c(sccs.size());
      for(size_t i=0;i<sccs.size();++i){</pre>
        for(size_t j=0;j<sccs[i].size();++j){</pre>
           result[sccs[i][j]]=c[i];
           c[idx[sccs[i][j]^1]]=!c[i];
        }
      return true;
    bool get(int x){return result[x];}
    inline int get_id(int x){return idx[x];}
    inline int count(){return sccs.size();}
} sat2:
```

## Lowbit Decomposition

```
class LowbitDecomp{
private:
  int time_, chain_, LOG_N;
  vector< vector< int > > G, fa;
vector< int > tl, tr, chain, chain_st;
  // chain_ : number of chain
  // tl, tr[ u ] : subtree interval in the seq. of u
  // chain_st[ u ] : head of the chain contains u // chian[ u ] : chain id of the chain u is on
  inline int lowbit( int x ) {
     return x & ( -x );
  void predfs( int u, int f ) {
     chain[ u ] = 0;
     for ( int v : G[ u ] ) {
       if ( v == f ) continue;
       predfs( v, u );
if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
          chain[ u ] = chain[ v ];
     if ( not chain[ u ] )
       chain[ u ] = chain_ ++;
  void dfschain( int u, int f ) {
    fa[ u ][ 0 ] = f;

for ( int i = 1 ; i < LOG_N ; ++ i )

fa[ u ][ i ] = fa[ fa[ u ][ i - 1 ] ][ i - 1 ];
     tl[ u ] = time_++
     if ( not chain_st[ chain[ u ] ] )
     chain_st[ chain[ u ] ] = u;
for ( int v : G[ u ] )
       if ( v != f and chain[ v ] == chain[ u ] )
     dfschain( v, u );
for ( int v : G[ u ] )
       if ( v != f and chain[ v ] != chain[ u ] )
          dfschain( v, u );
    tr[ u ] = time_;
  inline bool anc( int u, int v ) {
  return tl[ u ] <= tl[ v ] \</pre>
       and tr[ v ] <= tr[ u ];
public:
  inline int lca( int u, int v ) {
     if ( anc( u, v ) ) return u;
     for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
  if ( not anc( fa[ u ][ i ], v ) )
          u = fa[ u ][ i ];
     return fa[ u ][ 0 ];
```

```
void init( int n ) {
    n ++;
    for (LOG_N = 0 ; (1 << LOG_N) < n ; ++ LOG_N);
    fa.clear();
    fa.resize( n, vector< int >( LOG_N ) );
    G.clear(); G.resize( n );
    tl.clear(); tl.resize( n );
tr.clear(); tr.resize( n );
    chain.clear(); chain.resize( n );
    chain_st.clear(); chain_st.resize( n );
  void add_edge( int u , int v ) {
    // 1-base
    G[ u ].push_back( v );
    G[ v ].push_back( u );
  void decompose(){
    chain_ = 1;
    predfs( 1, 1 );
    time_{-} = 0;
    dfschain(1,1);
  PII get_inter( int u ) { return {tl[ u ], tr[ u ]}; }
  vector< PII > get_path( int u , int v ){
    vector< PII > res;
    int g = lca( u, v );
while ( chain[ u ] != chain[ g ] ) {
       int s = chain_st[ chain[ u ] ];
       res.emplace_back( tl[ s ], tl[ u ] + 1 );
       u = fa[ s ][ 0 ];
    res.emplace_back( tl[ g ], tl[ u ] + 1 );
while ( chain[ v ] != chain[ g ] ) {
  int s = chain_st[ chain[ v ] ];
       res.emplace_back( tl[ s ], tl[ v ] + 1 );
       v = fa[ s ][ 0 ];
    res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
    return res;
    /* res : list of intervals from u to v
     * ( note only nodes work, not edge )
     * usage :
     * vector< PII >& path = tree.get_path( u , v )
     * for( auto [ 1, \dot{r} ] : path ) {
         0-base [ 1, r )
} tree;
```

#### 3.6 MaxClique

```
// contain a self loop u to u, than u won't in clique
template < size_t MAXN >
class MaxClique{
private:
  using bits = bitset< MAXN >;
  bits popped, G[ MAXN ], ans;
  size_t deg[ MAXN ], deo[ MAXN ], n;
  void sort_by_degree() {
    popped.reset();
    for ( size_t i = 0 ; i < n ; ++ i )</pre>
    deg[ i ] = G[ i ].count();
for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;</pre>
         for ( size_t j = 0 ; j < n ; ++ j )</pre>
             if ( not popped[ j ] and deg[ j ] < mi )</pre>
                  mi = deg[id = j];
         popped[ deo[ i ] = id ] = 1;
         for( size_t u = G[ i ]._Find_first()
           u < n ; u = G[ i ]._Find_next( u ) )
             -- deg[ u ];
  void BK( bits R, bits P, bits X ) {
    if (R.count()+P.count() <= ans.count()) return;</pre>
    if ( not P.count() and not X.count() ) {
      if ( R.count() > ans.count() ) ans = R;
       return;
     /* greedily chosse max degree as pivot
```

```
bits cur = P | X; size_t pivot = 0, sz = 0;
for ( size_t u = cur._Find_first() ;
       u < n ; u = cur._Find_next( u ) )</pre>
         if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
     cur = P & ( \sim G[ pivot ]
     */ // or simply choose first
     bits cur = P & (~G[ ( P | X )._Find_first() ]);
     for ( size_t u = cur._Find_first() ;
       u < n ; u = cur._Find_next( u ) ) {
       if ( R[ u ] ) continue;
R[ u ] = 1;
       BK( R, P & G[ u ], X & G[ u ] );
       R[u] = P[u] = 0, X[u] = 1;
  }
public:
  void init( size_t n_ ) {
     n = n_{-}
     for ( size_t i = 0 ; i < n ; ++ i )
       G[ i ].reset();
     ans.reset();
  void add_edges( int u, bits S ) { G[ u ] = S; }
  void add_edge( int u, int v )
     G[u][v] = G[v][u] = 1;
  int solve() {
     sort_by_degree(); // or simply iota( deo... )
     for ( size_t i = 0 ; i < n ; ++ i )
     deg[ i ] = G[ i ].count();
bits pob, nob = 0; pob.set();
     for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
     for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t v = deo[ i ];</pre>
       bits tmp; tmp[ v ] = 1;
       BK( tmp, pob & G[ v ], nob & G[ v ] );
pob[ v ] = 0, nob[ v ] = 1;
     return static_cast< int >( ans.count() );
};
```

### 3.7 Virtural Tree

```
inline bool cmp(const int &i, const int &j) {
  return dfn[i] < dfn[j];</pre>
void build(int vectrices[], int k) {
  static int stk[MAX_N];
  sort(vectrices, vectrices + k, cmp);
  stk[sz++] = 0;
  for (int i = 0; i < k; ++i) {
  int u = vectrices[i], lca = LCA(u, stk[sz - 1]);</pre>
    if (lca == stk[sz - 1]) stk[sz++] = u;
    else {
      while (sz \ge 2 \&\& dep[stk[sz - 2]] \ge dep[lca]) {
        addEdge(stk[sz - 2], stk[sz - 1]);
        sz--;
      if (stk[sz - 1] != lca) {
        addEdge(lca, stk[--sz]);
        stk[sz++] = lca, vectrices[cnt++] = lca;
      stk[sz++] = u;
  for (int i = 0; i < sz - 1; ++i)
    addEdge(stk[i], stk[i + 1]);
```

# 3.8 Tree Hashing

```
uint64_t hsah( int u, int f ) {
    uint64_t r = 127;
    for ( int v : G[ u ] ) {
        if ( v == f ) continue;
        uint64_t hh = hsah( v, u );
        r = r + ( hh * hh ) % mod;
    }
    return r;
}
```

# 3.9 Minimum Mean Cycle

```
/* minimum mean cycle O(VE) */
struct MMC{
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
#define V 1021
#define inf 1e9
   struct Edge { int v,u; double c; };
   int n, m, prv[V][V], prve[V][V], vst[V];
   Edge e[E];
   vector<int> edgeID, cycle, rho;
   double d[V][V];
  void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
   void add_edge( int vi , int ui , double ci )
  { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
     for(int i=0; i<n; i++) d[0][i]=0;</pre>
     for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);</pre>
       for(int j=0; j<m; j++)</pre>
          int v = e[j].v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
            d[i+1][u] = d[i][v]+e[j].c;
            prv[i+1][u] = v;
            prve[i+1][u] = j;
         }
       }
    }
   double solve(){
     // returns inf if no cycle, mmc otherwise
     double mmc=inf;
     int st = -1
     bellman_ford();
     for(int i=0; i<n; i++) {</pre>
       double avg=-inf;
       for(int k=0; k<n; k++) {
  if(d[n][i]<inf-eps)</pre>
            avg=max(avg,(d[n][i]-d[k][i])/(n-k));
          else avg=max(avg,inf);
       if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
     FZ(vst);edgeID.clear();cycle.clear();rho.clear();
     for (int i=n; !vst[st]; st=prv[i--][st]) {
       vst[st]++;
       edgeID.PB(prve[i][st]);
       rho.PB(st);
     while (vst[st] != 2) {
       int v = rho.back(); rho.pop_back();
       cycle.PB(v);
       vst[v]++;
     reverse(ALL(edgeID));
     edgeID.resize(SZ(cycle));
     return mmc;
   }
} mmc;
```

#### 3.10 Mo's Algorithm on Tree

```
int n, q, nxt[ N ], to[ N ], hd[ N ];
struct Que{
  int u, v, id;
} que[ N ];
void init() {
  cin >> n >> q;
  for ( int i = 1 ; i < n ; ++ i ) {
    int u, v; cin >> u >> v;
    nxt[ i << 1 | 0 ] = hd[ u ];
    to[ i << 1 | 0 ] = v;
    hd[ u ] = i << 1 | 0;
    nxt[ i << 1 | 1 ] = hd[ v ];
    to[ i << 1 | 1 ] = u;
    hd[ v ] = i << 1 | 1;
}
for ( int i = 0 ; i < q ; ++ i ) {
    cin >> que[ i ].v; que[ i ].id = i;
```

```
}
int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
void dfs( int u, int f ) {
   dfn[ u ] = dfn_++; int saved_rbp = stk_;
  for ( int v_ = hd[ u ] ; v_ ; v_ = nxt[ v_ ] ) {
     if ( to[ v_ ] == f ) continue;
     dfs( to[ v_ ], u );
    if ( stk_ - saved_rbp < SQRT_N ) continue;
for ( ++ block_ ; stk_ != saved_rbp ; )
   block_id[ stk[ -- stk_ ] ] = block_;</pre>
  stk[ stk_ ++ ] = u;
bool inPath[ N ];
void Diff( int u ) {
  if ( inPath[ u ] ^= 1 )
    // remove this edge
  else
    // add this edge
void traverse( int& origin_u, int u ) {
  for ( int g = lca( origin_u, u )
     origin_u != g ; origin_u = parent_of[ origin_u ] )
       Diff( origin_u );
  for (int v = u; v != origin_u; v = parent_of[v])
    Diff( v );
  origin_u = u;
void solve() {
  while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
  sort( que, que + q, [](const Que& x, const Que& y) {
  return tie( block_id[ x.u ], dfn[ x.v ] )
               < tie( block_id[ y.u ], dfn[ y.v ] );
  } );
  int U = 1, V = 1;
  for ( int i = 0 ; i < q ; ++ i ) {
    pass( U, que[ i ].u );
pass( V, que[ i ].v );
     // we could get our answer of que[ i ].id
Method 2:
dfs u:
  push u
  iterate subtree
  push u
Let P = LCA(u, v), and St(u) <= St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
```

#### 3.11 Minimum Steiner Tree

```
// Minimum Steiner Tree
// 0(V 3^T + V^2 2^T)
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
  int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
  void init( int _n ){
    n = _n;
     for( int i = 0 ; i < n ; i ++ ){</pre>
       for( int j = 0 ; j < n ; j ++ )
  dst[ i ][ j ] = INF;
dst[ i ][ i ] = 0;</pre>
    }
  void add_edge( int ui , int vi , int wi ){
    dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
  void shortest_path(){
     for( int k = 0 ; k < n ; k ++ )</pre>
       for( int i = 0 ; i < n ; i ++ )</pre>
          for( int j = 0 ; j < n ; j ++ )
  dst[ i ][ j ] = min( dst[ i ][ j ],</pre>
                     dst[ i ][ k ] + dst[ k ][ j ] );
```

```
int solve( const vector<int>& ter ){
     int t = (int)ter.size();
     for( int i = 0 ; i < ( 1 << t ) ; i ++ )
    for( int j = 0 ; j < n ; j ++ )
dp[ i ][ j ] = INF;
for( int i = 0 ; i < n ; i ++ )
       dp[0][i] = 0;
     for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
       if( msk == ( msk & (-msk) ) ){
         int who = __lg( msk );
for( int i = 0 ; i < n ; i ++ )
            dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
         continue;
       for( int i = 0 ; i < n ; i ++ )
         for( int submsk = ( msk - 1 ) & msk ; submsk ;
    submsk = ( submsk - 1 ) & msk )
              dp[ msk ][ i ] = min( dp[ msk ][ i ],
                                 dp[ submsk ][ i ] +
                                  dp[ msk ^ submsk ][ i ] );
       for( int i = 0 ; i < n ; i ++ ){</pre>
         tdst[ i ] = INF;
for( int j = 0 ; j < n ; j ++ )
            tdst[ i ] = min( tdst[ i ],
                         dp[ msk ][ j ] + dst[ j ][ i ] );
       for( int i = 0 ; i < n ; i ++ )</pre>
          dp[ msk ][ i ] = tdst[ i ];
     int ans = INF;
     for( int i = 0 ; i < n ; i ++ )</pre>
       ans = min( ans , dp[ ( 1 << t ) - 1 ][ i ] );
     return ans;
} solver;
```

## 3.12 Directed Minimum Spanning Tree

```
template <typename T> struct DMST {
  T g[maxn][maxn], fw[maxn];
  int n, fr[maxn];
  bool vis[maxn], inc[maxn];
  void clear() {
    for(int i = 0; i < maxn; ++i) {</pre>
      for(int j = 0; j < maxn; ++j) g[i][j] = inf;</pre>
      vis[i] = inc[i] = false;
    }
  void addEdge(int u,int v,T w){g[u][v]=min(g[u][v],w)
  T operator()(int root, int _n) {
    n = _n; T ans = 0;
    if (dfs(root) != n) return -1;
    while (true) {
      for(int i = 1;i <= n;++i) fw[i] = inf, fr[i] = i;</pre>
      for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
        for (int j = 1; j <= n; ++j) {</pre>
          if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
            fw[i] = g[j][i]; fr[i] = j;
        }
      int x = -1;
      for(int i = 1;i <= n;++i)if(i != root && !inc[i])</pre>
         int j = i, c = 0;
        while(j!=root && fr[j]!=i && c<=n) ++c, j=fr[j</pre>
             ];
        if (j == root || c > n) continue;
        else { x = i; break; }
      if (!~x) {
         for (int i = 1; i <= n; ++i)</pre>
          if (i != root && !inc[i]) ans += fw[i];
        return ans;
      int y = x;
      for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
      do {
```

```
ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true
;
} while (y != x);
inc[x] = false;
for (int k = 1; k <= n; ++k) if (vis[k]) {
    for (int j = 1; j <= n; ++j) if (!vis[j]) {
        if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
        if (g[j][k] < inf && g[j][k]-fw[k] < g[j][x])
        g[j][x] = g[j][k] - fw[k];
}

}
return ans;
}
int dfs(int now) {
  int r = 1; vis[now] = true;
  for (int i = 1; i <= n; ++i)
    if (g[now][i] < inf && !vis[i]) r += dfs(i);
  return r;
}
};</pre>
```

#### 3.13 Dominator Tree

```
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
  // vertices are numbered from 0 to n-1
 fill(dfn, dfn + n, -1);fill(rev, rev + n, -1);
  fill(fa, fa + n, -1); fill(val, val + n, -1);
  fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
 fill(dom, dom + n, -1); tk = 0;
 for (int i = 0; i < n; ++i) {
   g[i].clear(); r[i].clear(); rdom[i].clear();
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
 rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
  for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
    r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0)
 if (fa[x] == x) return c ? -1 : x;
 int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
  if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
  fa[x] = p;
  return c ? p : val[x];
vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
// p[i] = -2 if i is unreachable from s
 dfs(s);
  for (int i = tk - 1; i \ge 0; --i) {
    for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)])
    if (i) rdom[sdom[i]].push_back(i);
    for (int &u : rdom[i]) {
      int p = find(u);
      if (sdom[p] == i) dom[u] = i;
      else dom[u] = p;
    if (i) merge(i, rp[i]);
  vector<int> p(n, -2); p[s] = -1;
  for (int i = 1; i < tk; ++i)
    if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
  for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
  return p;
}}
```

# 4 Matching & Flow

## 4.1 Kuhn Munkres

```
class KM {
private:
  static constexpr lld INF = 1LL << 60;</pre>
  vector<lld> hl,hr,slk;
  vector<int> fl,fr,pre,qu;
vector<vector<lld>> w;
  vector<bool> v1,vr;
  int n, ql, qr;
bool check(int x) {
     if (v1[x] = true, f1[x] != -1)
    return vr[qu[qr++] = f1[x]] = true;
while (x != -1) swap(x, fr[f1[x] = pre[x]]);
     return false;
  void bfs(int s) {
     fill(slk.begin(), slk.end(), INF);
    fill(vl.begin(), vl.end(), false);
fill(vr.begin(), vr.end(), false);
     ql = qr = 0;
     qu[qr++] = s;
     vr[s] = true;
    while (true) {
       11d d;
       while (ql < qr) {
  for (int x = 0, y = qu[ql++]; x < n; ++x) +</pre>
            if(!v1[x]&&slk[x]>=(d=h1[x]+hr[y]-w[x][y])){
              if (pre[x] = y, d) slk[x] = d;
else if (!check(x)) return;
           }
         }
       d = INF;
       for (int x = 0; x < n; ++x)
         if (!v1[x] \&\& d > s1k[x]) d = s1k[x];
       for (int x = 0; x < n; ++x) {
         if (v1[x]) h1[x] += d;
         else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
       for (int x = 0; x < n; ++x)
         if (!v1[x] && !slk[x] && !check(x)) return;
public:
  void init( int n_ ) {
    n = n_; qu.resize(n);
fl.clear(); fl.resize(n, -1);
     fr.clear(); fr.resize(n, -1);
    hr.clear(); hr.resize(n); hl.resize(n);
     w.clear(); w.resize(n, vector<lld>(n));
     slk.resize(n); pre.resize(n);
     vl.resize(n); vr.resize(n);
  void set_edge( int u, int v, lld x ) {w[u][v] = x;}
  1ld solve() {
     for (int i = 0; i < n; ++i)
      hl[i] = *max_element(w[i].begin(), w[i].end());
     for (int i = 0; i < n; ++i) bfs(i);</pre>
     11d res = 0;
     for (int i = 0; i < n; ++i) res += w[i][f1[i]];</pre>
     return res;
} km;
```

#### 4.2 Bipartite Matching

```
class BipartiteMatching{
private:
    vector<int> X[N], Y[N];
    int fX[N], fY[N], n;
    bitset<N> walked;
    bool dfs(int x){
        for(auto i:X[x]){
            if(walked[i])continue;
            walked[i]=1;
            if(fY[i]==-1||dfs(fY[i])){
                 fY[i]=x;fX[x]=i;
```

```
return 1:
       }
    return 0:
public:
  void init(int _n){
    n=_n; walked.reset();
    for(int i=0;i<n;i++){</pre>
       X[i].clear();Y[i].clear();
       fX[i]=fY[i]=-1;
    }
  void add_edge(int x, int y){
  X[x].push_back(y); Y[y].push_back(y);
  int solve(){
    int cnt = 0;
    for(int i=0;i<n;i++){</pre>
       walked.reset();
       if(dfs(i)) cnt++;
    // return how many pair matched
    return cnt;
};
```

## 4.3 General Graph Matching

```
const int N = 514, E = (2e5) * 2;
struct Graph{
  int to[E],bro[E],head[N],e;
  int lnk[N], vis[N], stp, n;
  void init( int _n ){
  stp = 0; e = 1; n = _n;
    for( int i = 0 ; i <= n ; i ++ )</pre>
      head[i] = lnk[i] = vis[i] = 0;
  void add_edge(int u,int v){
    // 1-base
    to[e]=v,bro[e]=head[u],head[u]=e++;
    to[e]=u,bro[e]=head[v],head[v]=e++;
  bool dfs(int x){
    vis[x]=stp;
    for(int i=head[x];i;i=bro[i]){
      int v=to[i];
      if(!lnk[v]){
        lnk[x]=v, lnk[v]=x;
         return true
      }else if(vis[lnk[v]]<stp){</pre>
        int w=lnk[v]
         lnk[x]=v, lnk[v]=x, lnk[w]=0;
        if(dfs(w)) return true
        lnk[w]=v, lnk[v]=w, lnk[x]=0;
      }
    }
    return false;
  int solve(){
    int ans = 0;
    for(int i=1;i<=n;i++)</pre>
      if(not lnk[i]){
        stp++; ans += dfs(i);
    return ans;
} graph;
```

# 4.4 Minimum Weight Matching (Clique version)

```
struct Graph {
  // 0-base (Perfect Match)
  int n, edge[MXN][MXN];
  int match[MXN], dis[MXN], onstk[MXN];
  vector<int> stk;
  void init(int _n) {
    n = _n;
    for (int i=0; i<n; i++)</pre>
```

```
for (int j=0; j<n; j++)</pre>
        edge[i][j] = 0;
  void set_edge(int u, int v, int w) {
    edge[u][v] = edge[v][u] = w;
  bool SPFA(int u){
    if (onstk[u]) return true;
    stk.PB(u);
    onstk[u] = 1
    for (int v=0; v<n; v++){
      if (u != v && match[u] != v && !onstk[v]){
        int m = match[v];
        if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
          dis[m] = dis[u] - edge[v][m] + edge[u][v];
          onstk[v] = 1;
          stk.PB(v);
          if (SPFA(m)) return true;
          stk.pop_back();
          onstk[v] = 0;
      }
    onstk[u] = 0;
    stk.pop_back();
    return false;
  int solve() {
    // find a match
    for (int i=0; i<n; i+=2){</pre>
      match[i] = i+1;
      match[i+1] = i;
    while (true){
      int found = 0;
      for (int i=0; i<n; i++)</pre>
        dis[i] = onstk[i] = 0;
      for (int i=0; i<n; i++){</pre>
        stk.clear()
        if (!onstk[i] && SPFA(i)){
          found = 1
          while (SZ(stk)>=2){
            int u = stk.back(); stk.pop_back();
             int v = stk.back(); stk.pop_back();
            match[u] = v;
            match[v] = u;
          }
        }
      if (!found) break;
    }
    int ret = 0;
    for (int i=0; i<n; i++)
      ret += edge[i][match[i]];
    return ret>>1;
} graph;
```

#### 4.5 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source  ${\it S}$  and sink  ${\it T}$ .
  - 2. For each edge (x,y,l,u), connect  $x \to y$  with capacity u-l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v\to T$  with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer
    - from s to t is the answer.

      To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$  , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching  ${\cal M}$  on bipartite graph (X,Y)

```
1. Redirect every edge: y \to x if (x,y) \in M, x \to y other-
    2. DFS from unmatched vertices in X.
3. x \in X is chosen iff x is unvisited.
4. y \in Y is chosen iff y is visited.

    Minimum cost cyclic flow

     1. Consruct super source S and sink T 2. For each edge (x,y,c), connect x \to y with (cost,cap) =
         (c,1) if c>0, otherwise connect y \to x with (cost, cap) =
     3. For each edge with c\,<\,0 , sum these cost as K , then
     increase d(y) by 1, decrease d(x) by 1 4. For each vertex v with d(v)>0, connect S\to v with
         (cost, cap) = (0, d(v))
     5. For each vertex v with d(v) < 0, connect v 	o T with
     (cost, cap) = (0, -d(v)) 6. Flow from S to T, the answer is the cost of the flow
         C + K
· Maximum density induced subgraph
     1. Binary search on answer, suppose we're checking answer
     2. Construct a max flow model, let K be the sum of all
        weights
     3. Connect source s \to v, v \in G with capacity K
     4. For each edge (u,v,w) in G, connect u 	o v and v 	o u
         with capacity \hat{w}
     5. For v\in G, connect it with sink v\to t with capacity K+2T-(\sum_{e\in E(v)}w(e))-2w(v)
     6. T is a valid answer if the maximum flow f < K|V|

    Minimum weight edge cover

     1. For each v \in V create a copy v', and connect u' 	o v'
    with weight w(u,v). 
 2. Connect v \to v' with weight 2\mu(v), where \mu(v) is the cost
         of the cheapest edge incident to v\,.
     3. Find the minimum weight perfect matching on G'.
```

#### · Project selection problem

- 1. If  $p_v > 0$ , create edge (s, v) with capacity  $p_v$ ; otherwise,
- create edge (v,t) with capacity  $-p_v$ . 2. Create edge (u,v) with capacity w with w being the cost
- of choosing u without choosing v. 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge  $\left(x,t\right)$  with capacity  $c_{x}$  and create edge  $\left(s,y\right)$ with capacity  $c_y$ . 2. Create edge (x,y) with capacity  $c_{xy}$ . 3. Create edge (x,y) and edge (x',y') with capacity  $c_{xyx'y'}$ .

#### 4.6 Dinic

```
class Dinic{
private:
 using CapT = int64_t;
  struct Edge{
    int to, rev;
    CapT cap:
  int n, st, ed;
  vector<vector<Edge>> G;
  vector<int> lv, idx;
 bool BFS(){
    fill(lv.begin(), lv.end(), -1);
    queue<int> bfs;
    bfs.push(st);
    lv[st] = 0;
    while(!bfs.empty()){
      int u = bfs.front(); bfs.pop();
      for(auto e: G[u]){
        if(e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
        lv[e.to] = lv[u] + 1;
        bfs.push(e.to);
      }
    return (lv[ed]!=-1);
  CapT DFS(int u, CapT f){
    if(u == ed) return f;
    CapT ret = 0;
    for(int& i = idx[u]; i < (int)G[u].size(); ++i){</pre>
      auto& e = G[u][i]
      if(e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
```

```
CapT nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
       if(f == 0) return ret;
     if(ret == 0) lv[u] = -1;
     return ret;
public:
  void init(int n_, int st_, int ed_){
     n = n_, st = st_, ed = ed_;
G.resize(n); lv.resize(n);
     fill(G.begin(), G.end(), vector<Edge>());
  void add_edge(int u, int v, CapT c){
     G[u].push_back({v, (int)G[v].size(), c});
G[v].push_back({u, ((int)G[u].size())-1, 0});
  CapT max_flow(){
     CapT ret = 0;
     while(BFS()){
       idx.assign(n, 0);
       CapT f = DFS(st, numeric_limits<CapT>::max());
       ret += f;
       if(f == 0) break;
     return ret;
} flow;
```

#### 4.7 Minimum Cost Maximum Flow

```
class MiniCostMaxiFlow{
  using CapT = int;
  using WeiT = int64_t;
  using PCW = pair<CapT,WeiT>;
  static constexpr CapT INF_CAP = 1 << 30;</pre>
  static constexpr WeiT INF_WEI = 1LL<<60;</pre>
private:
  struct Edge{
    int to, back;
    WeiT wei;
    CapT cap
    Edge() {}
    Edge(int a,int b,WeiT c,CapT d):
      to(a), back(b), wei(c), cap(d)
    {}
  int ori, edd;
  vector<vector<Edge>> G;
  vector<int> fa, wh;
  vector<bool> inq;
  vector<WeiT> dis;
  PCW SPFA(){
    fill(inq.begin(),inq.end(),false);
    fill(dis.begin(), dis.end(), INF_WEI);
    queue<int> qq; qq.push(ori);
    dis[ori]=0;
    while(!qq.empty()){
      int u=qq.front();qq.pop();
      inq[u] = 0;
      for(int i=0;i<SZ(G[u]);++i){</pre>
        Edge e=G[u][i];
        int v=e.to:
        WeiT d=e.wei;
        if(e.cap<=0||dis[v]<=dis[u]+d)</pre>
          continue
        dis[v]=dis[u]+d;
        fa[v]=u,wh[v]=i;
        if(inq[v]) continue;
        qq.push(v);
        inq[v]=1;
      }
    if(dis[edd]==INF_WEI)
      return {-1,-1};
    CapT mw=INF_CAP;
    for(int i=edd;i!=ori;i=fa[i])
      mw=min(mw,G[fa[i]][wh[i]].cap);
    for (int i=edd;i!=ori;i=fa[i]){
      auto &eg=G[fa[i]][wh[i]];
```

```
eq.cap-=mw:
      G[eg.to][eg.back].cap+=mw;
    return {mw,dis[edd]};
public:
 void init(int a,int b,int n){
    ori=a,edd=b;
    G.clear();G.resize(n);
    fa.resize(n);wh.resize(n);
    inq.resize(n); dis.resize(n);
  void add_edge(int st,int ed,WeiT w,CapT c){
    G[st].emplace_back(ed,SZ(G[ed]),w,c);
    G[ed].emplace_back(st,SZ(G[st])-1,-w,0);
  PCW solve(){
    /* might modify to
    cc += ret.first * ret.second
    ww += ret.first * ret.second
    CapT cc=0; WeiT ww=0;
    while(true)
      PCW ret=SPFA();
      if(ret.first==-1) break;
      cc+=ret.first;
      ww+=ret.second;
    }
    return {cc,ww};
} mcmf;
```

#### 4.8 Global Min-Cut

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
    w[x][y] += c;
    w[y][x] += c;
pair<int, int> phase(int n) {
    memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
    int s = -1, t = -1;
    while (true) {
         int c = -1;
         for (int i = 0; i < n; ++i) {</pre>
              if (del[i] || v[i]) continue;
              if (c == -1 \mid | g[i] > g[c]) c = i;
         if (c == -1) break;
         v[c] = true;
         s = t, t = c;
         for (int i = 0; i < n; ++i) {
    if (del[i] || v[i]) continue;</pre>
              g[i] += w[c][i];
     return make_pair(s, t);
int mincut(int n) {
    int cut = 1e9;
    memset(del, false, sizeof(del));
     for (int i = 0; i < n - 1; ++i) {
         int s, t; tie(s, t) = phase(n);
         del[t] = true;
         cut = min(cut, g[t]);
for (int j = 0; j < n; ++j) {
    w[s][j] += w[t][j];</pre>
              w[j][s] += w[j][t];
     return cut;
}
```

## 5 Math

#### 5.1 Prime Table

```
\begin{array}{c} 1002939109, 1020288887, 1028798297, 1038684299, \\ 1041211027, 1051762951, 1058585963, 1063020809, \\ 1147930723, 1172520109, 1183835981, 1187659051, \\ 1241251303, 1247184097, 1255940849, 1272759031, \\ 1287027493, 1288511629, 1294632499, 1312650799, \\ 1868732623, 1884198443, 1884616807, 1885059541, \\ 1909942399, 1914471137, 1923951707, 1925453197, \\ 1979612177, 1980446837, 1989761941, 2007826547, \\ 2008033571, 2011186739, 2039465081, 2039728567, \\ 2093735719, 2116097521, 2123852629, 2140170259, \\ 3148478261, 3153064147, 3176351071, 3187523093, \\ 3196772239, 3201312913, 3203063977, 3204840059, \\ 3210224309, 3213032591, 3217689851, 3218469083, \\ 3219857533, 3231880427, 3235951699, 3273767923, \\ 3276188869, 3277183181, 3282463507, 3285553889, \\ 3319309027, 3327005333, 3327574903, 3341387953, \\ 3373293941, 3380077549, 3380892997, 3381118801. \end{aligned}
```

# 5.2 $\lfloor \frac{n}{i} \rfloor$ Enumeration

 $T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_{i+1} \rfloor} \rfloor}$ 

## 5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b==(x/y)*a;
}
```

#### 5.4 Pollard Rho

```
// does not work when n is prime
 // return any non-trivial factor
llu pollard_rho(llu n){
  static auto f=[](llu x,llu k,llu m){
    return add(k,mul(x,x,m),m);
  if (!(n&1)) return 2;
  mt19937 rnd(120821011);
  while(true){
    1lu y=2, yy=y, x=rnd()%n, t=1;
    for(llu sz=2;t==1;sz<<=1) {</pre>
      for(llu i=0;i<sz;++i){</pre>
        if(t!=1)break;
        yy=f(yy,x,n);
        t=gcd(yy>y?yy-y:y-yy,n);
      y=yy;
    if(t!=1&&t!=n) return t;
}
```

## 5.5 Pi Count (Linear Sieve)

```
static constexpr int N = 1000000 + 5;
11d pi[N];
vector<int> primes;
bool sieved[N];
1ld cube_root(lld x){
  1ld s=cbrt(x-static_cast<long double>(0.1));
  while(s*s*s <= x) ++s;
  return s-1;
1ld square_root(1ld x){
  1ld s=sqrt(x-static_cast<long double>(0.1));
  while(s*s <= x) ++s;
  return s-1:
void init(){
  primes.reserve(N);
  primes.push_back(1);
  for(int i=2;i<N;i++) {</pre>
    if(!sieved[i]) primes.push_back(i);
    pi[i] = !sieved[i] + pi[i-1];
    for(int p: primes) if(p > 1) {
      if(p * i >= N) break;
      sieved[p * i] = true;
      if(p % i == 0) break;
```

```
}
11d phi(11d m, 11d n) {
  static constexpr int MM = 80000, NN = 500;
  static lld val[MM][NN];
  if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
  if(n == 0) return m;
  if(primes[n] >= m) return 1;
lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
  if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
  return ret:
11d pi_count(11d);
11d P2(11d m, 11d n) {
  11d sm = square_root(m), ret = 0;
  for(lld i = n+1;primes[i]<=sm;i++)</pre>
    ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
  return ret;
11d pi_count(11d m) {
  if(m < N) return pi[m];</pre>
  11d n = pi_count(cube_root(m));
  return phi(m, n) + n - 1 - P2(m, n);
```

## 5.6 Range Sieve

#### 5.7 Miller Rabin

```
bool isprime(llu x){
 static llu magic[]={2,325,9375,28178,\
                     450775,9780504,1795265022};
  static auto witn=[](llu a,llu u,llu n,int t){
    a = mpow(a,u,n);
    if (!a)return 0;
   while(t--){
      1lu a2=mul(a,a,n);
      if(a2==1 && a!=1 && a!=n-1)
        return 1;
      a = a2;
    return a!=1;
 if(x<2)return 0;</pre>
  if(!(x&1))return x==2;
 llu x1=x-1; int t=0;
 while(!(x1&1))x1>>=1,t++;
  for(llu m:magic)if(witn(m,x1,x,t))return 0;
  return 1:
```

#### 5.8 Inverse Element

```
// x's inverse mod k
long long GetInv(long long x, long long k){
  // k is prime: euler_(k)=k-1
  return qPow(x, euler_phi(k)-1);
}
// if you need [1, x] (most use: [1, k-1]
void solve(int x, long long k){
```

```
inv[1] = 1;
for(int i=2;i<x;i++)
  inv[i] = ((long long)(k - k/i) * inv[k % i]) % k;
}</pre>
```

#### 5.9 Euler Phi Function

```
extended euler:
   a^b mod p
   if gcd(a, p)==1: a^{(b\%phi(p))}
   elif b < phi(p): a^b mod p
   else a^(b%phi(p) + phi(p))
lld euler_phi(int x){
  11d r=1;
  for(int i=2;i*i<=x;++i){</pre>
    if(x%i==0){
      x/=i; r*=(i-1);
      while(x%i==0){
        x/=i; r*=i;
    }
  if(x>1) r*=x-1;
  return r;
vector<int> primes;
bool notprime[N];
1ld phi[N];
void euler_sieve(int n){
  for(int i=2;i<n;i++){</pre>
    if(!notprime[i]){
      primes.push_back(i); phi[i] = i-1;
    for(auto j: primes){
      if(i*j >= n) break;
      notprime[i*j] = true;
      phi[i*j] = phi[i] * phi[j];
      if(i % j == 0){
        phi[i*j] = phi[i] * j;
        break;
  }
}
```

#### 5.10 Gauss Elimination

```
void gauss(vector<vector<double>> &d) {
   int n = d.size(), m = d[0].size();
   for (int i = 0; i < m; ++i) {
      int p = -1;
      for (int j = i; j < n; ++j) {
        if (fabs(d[j][i]) < eps) continue;
        if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
      }
      if (p == -1) continue;
      for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
      for (int j = 0; j < n; ++j) {
        if (i == j) continue;
        double z = d[j][i] / d[i][i];
      for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
      }
    }
}</pre>
```

## 5.11 Fast Fourier Transform

```
polynomial multiply:
   DFT(a, len); DFT(b, len);
   for(int i=0;i<len;i++) c[i] = a[i]*b[i];
   iDFT(c, len);
   (len must be 2^k and >= 2*(max(a, b)))
   Hand written Cplx would be 2x faster
   */
Cplx omega[2][N];
void init_omega(int n) {
   static constexpr llf PI=acos(-1);
```

```
const llf arg=(PI+PI)/n;
  for(int i=0;i<n;++i)</pre>
    omega[0][i]={cos(arg*i),sin(arg*i)};
  for(int i=0;i<n;++i)</pre>
    omega[1][i]=conj(omega[0][i]);
void tran(Cplx arr[],int n,Cplx omg[]) {
  for(int i=0, j=0;i<n;++i){</pre>
    if(i>j)swap(arr[i],arr[j]);
    for(int l=n>>1;(j^=1)<1;l>>=1);
  for (int l=2;l<=n;l<<=1){</pre>
    int m=1>>1;
    for(auto p=arr;p!=arr+n;p+=1){
      for(int i=0;i<m;++i){</pre>
        Cplx t=omg[n/1*i]*p[m+i];
        p[m+i]=p[i]-t; p[i]+=t;
 }
void DFT(Cplx arr[],int n){tran(arr,n,omega[0]);}
void iDFT(Cplx arr[],int n){
  tran(arr,n,omega[1]);
  for(int i=0;i<n;++i) arr[i]/=n;</pre>
```

## 5.12 High Speed Linear Recurrence

```
#define mod 998244353
const int N=1000010;
int n,k,m,f[N],h[N],a[N],b[N],ib[N];
int pw(int x,int y){
  int re=1;
  if(y<0)y+=mod-1;
  while(y)
    if(y&1)re=(11)re*x%mod;
    y>=1; x=(11)x*x%mod;
  return re;
void inc(int&x,int y){x+=y;if(x>=mod)x-=mod;}
namespace poly{
  const int G=3;
  int rev[N],L;
  void ntt(int*A,int len,int f){
    for(L=0;(1<<L)<len;++L);
    for(int i=0;i<len;++i){</pre>
      rev[i]=(rev[i>>1]>>1)|((i&1)<<(L-1));
      if(i<rev[i])swap(A[i],A[rev[i]]);</pre>
    for(int i=1;i<len;i<<=1){</pre>
      int wn=pw(G, f*(mod-1)/(i<<1));</pre>
      for(int j=0;j<len;j+=i<<1){</pre>
        for(int k=0;k<i;++k,w=(11)w*wn%mod){</pre>
           int x=A[j+k], y=(11)w*A[j+k+i]%mod;
           A[j+k]=(x+y)\mbox{mod}, A[j+k+i]=(x-y+mod)\mbox{mod};
        }
      }
    if(!~f){
      int iv=pw(len,mod-2);
      for(int i=0;i<len;++i)A[i]=(11)A[i]*iv%mod;</pre>
  void cls(int*A,int l,int r){
    for(int i=1;i<r;++i)A[i]=0;}</pre>
  void cpy(int*A,int*B,int 1){
    for(int i=0;i<1;++i)A[i]=B[i];}</pre>
  void inv(int*A,int*B,int 1){
    if(l==1){B[0]=pw(A[0],mod-2);return;}
    static int t[N];
    int len=l<<1;
    inv(A,B,1>>1);
    cpy(t,A,1);cls(t,1,len);
    ntt(t,len,1);ntt(B,len,1);
    for(int i=0;i<len;++i)</pre>
      B[i]=(11)B[i]*(2-(11)t[i]*B[i]*mod+mod)*mod;
    ntt(B,len,-1);cls(B,l,len);
```

```
void pmod(int*A)
    static int t[N];
    int l=k+1,len=1;while(len<=(k<<1))len<<=1;</pre>
    cpy(t, A, (k<<1)+1);
    reverse(t,t+(k<<1)+1);
    cls(t,1,len);
    ntt(t,len,1)
    for(int i=0;i<len;++i)t[i]=(11)t[i]*ib[i]%mod;</pre>
    ntt(t,len,-1);
    cls(t,1,len)
    reverse(t,t+1);
    ntt(t,len,1)
    for(int i=0;i<len;++i)t[i]=(11)t[i]*b[i]%mod;</pre>
    ntt(t,len,-1);
    cls(t,l,len);
    for(int i=0;i<k;++i)A[i]=(A[i]-t[i]+mod)%mod;</pre>
    cls(A,k,len);
  void pow(int*A,int n){
    if(n==1) {cls(A,0,k+1);A[1]=1;return;}
    pow(A, n>>1);
    int len=1; while(len<=(k<<1))len<<=1;</pre>
    ntt(A,len,1);
    for(int i=0;i<len;++i)A[i]=(11)A[i]*A[i]%mod;</pre>
    ntt(A,len,-1);
    pmod(A)
    if(n&1){
      for(int i=k;i;--i)A[i]=A[i-1];A[0]=0;
      pmod(A);
  }
int main(){
  n=rd();k=rd();
  for(int i=1;i<=k;++i)f[i]=(mod+rd())%mod;</pre>
  for(int i=0;i<k;++i)h[i]=(mod+rd())%mod;</pre>
  for(int i=a[k]=b[k]=1;i<=k;++i)</pre>
    a[k-i]=b[k-i]=(mod-f[i])%mod;
  int len=1; while(len<=(k<<1))len<<=1;</pre>
  reverse(a,a+k+1);
  poly::inv(a,ib,len);
  poly::cls(ib,k+1,len);
  poly::ntt(b,len,1);
  poly::ntt(ib,len,1);
  poly::pow(a,n);
  int ans=0;
  for(int i=0;i<k;++i)inc(ans,(ll)a[i]*h[i]%mod);</pre>
  printf("%d\n",ans);
  return 0;
}
```

#### 5.13 Chinese Remainder

```
1ld crt(lld ans[], lld pri[], int n){
  lld M = 1, ret = 0;
  for(int i=0;i<n;i++) M *= pri[i];</pre>
  for(int i=0;i<n;i++){</pre>
    1ld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
    ret += (ans[i]*(M/pri[i])%M * iv)%M;
    ret %= M:
  }
  return ret;
}
/*
Another:
x = a1 \% m1
x = a2 \% m2
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
0 <= x < lcm(m1, m2)
*/
```

## 5.14 Berlekamp Massey

```
// x: 1-base, p[]: 0-base
template<size_t N>
vector<llf> BM(llf x[N], size_t n){
    size_t f[N]={0}, t=0;llf d[N];
    vector<llf> p[N];
```

```
for(size_t i=1,b=0;i<=n;++i) {
    for(size_t j=0;j<p[t].size();++j)
        d[i]+=x[i-j-1]*p[t][j];
    if(abs(d[i]-=x[i])<=EPS)continue;
    f[t]=i;if(!t){p[++t].resize(i);continue;}
    vector<llf>    cur(i-f[b]-1);
    llf k=-d[i]/d[f[b]];cur.PB(-k);
    for(size_t j=0;j<p[b].size();j++)
        cur.PB(p[b][j]*k);
    if(cur.size()<p[t].size())cur.resize(p[t].size());
    for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];
    if(i-f[b]+p[b].size()>=p[t].size()) b=t;
    p[++t]=cur;
}
return p[t];
}
```

#### 5.15 NTT

```
// Remember coefficient are mod P
/* p=a*2^n+1
                                       root
   n
   16
        65536
                     65537
                                       3
   20
        1048576
                     7340033
                                       3 */
// (must be 2<sup>k</sup>)
template<LL P, LL root, int MAXN>
struct NTT{
  static LL bigmod(LL a, LL b) {
    LL res = 1;
    for (LL bs = a; b; b >>= 1, bs = (bs * bs) % P)
      if(b&1) res=(res*bs)%P;
    return res;
 static LL inv(LL a, LL b) {
    if(a==1)return 1;
    return (((LL)(a-inv(b%a,a))*b+1)/a)%b;
 LL omega[MAXN+1];
 NTT() {
    omega[0] = 1;
    LL r = bigmod(root, (P-1)/MAXN);
    for (int i=1; i<=MAXN; i++)</pre>
      omega[i] = (omega[i-1]*r)%P;
  // n must be 2^k
  void tran(int n, LL a[], bool inv_ntt=false){
    int basic = MAXN / n , theta = basic;
for (int m = n; m >= 2; m >>= 1) {
      int mh = m >> 1;
      for (int i = 0; i < mh; i++)</pre>
        LL w = omega[i*theta%MAXN];
        for (int j = i; j < n; j += m) {
          int k = j + mh;
LL x = a[j] - a[k];
           if (x < 0) x += P;
          a[j] += a[k];
          if (a[j] > P) a[j] -= P;
          a[k] = (w * x) % P;
        }
      theta = (theta * 2) % MAXN;
    for (int j = 1; j < n - 1; j++) {
      for (int k = n >> 1; k > (i ^= k); k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
    if (inv_ntt) {
      LL ni = inv(n,P);
      reverse( a+1 , a+n );
      for (i = 0; i < n; i++)
        a[i] = (a[i] * ni) % P;
const LL P=2013265921, root=31;
const int MAXN=4194304;
NTT<P, root, MAXN> ntt;
```

## 5.16 Polynomial Operations

```
using VI = vector<int>;
Poly Inverse(Poly f) {
  int n = f.size()
  Poly q(1, fpow(f[0], kMod - 2));
  for (int s = 2;; s <<= 1)
    if (f.size() < s) f.resize(s);</pre>
    Poly fv(f.begin(), f.begin() + s);
    Poly fq(q.begin(), q.end());
    fv.resize(s + s); fq.resize(s + s);
    ntt::Transform(fv, s + s);
    ntt::Transform(fq, s + s);
    for (int i = 0; i < s + s; ++i)
      fv[i] = 1LL * fv[i] * fq[i]%kMod * fq[i]%kMod;
    ntt::InverseTransform(fv, s + s);
    Poly res(s);
    for (int i = 0; i < s; ++i) {
      res[i] = kMod - fv[i];
      if (i < (s >> 1))
        int v = 2 * q[i] % kMod;
         (res[i] += v) >= kMod ? res[i] -= kMod : 0;
      }
    q = res;
    if (s >= n) break;
  q.resize(n);
  return q;
Poly Divide(const Poly &a, const Poly &b) {
  int n = a.size(), m = b.size(), k = 2;
while (k < n - m + 1) k <<= 1;</pre>
  Poly ra(k), rb(k);
  for (int i = 0; i < min(n, k); ++i) ra[i] = a[n-1-i];
for (int i = 0; i < min(m, k); ++i) rb[i] = b[m-1-i];</pre>
  auto rbi = Inverse(rb);
  auto res = Multiply(rbi, ra);
  res.resize(n - m + 1);
  reverse(res.begin(), res.end());
  return res;
Poly Modulo(const Poly &a, const Poly &b) {
  if (a.size() < b.size()) return a;</pre>
  auto dv = Multiply(Divide(a, b), b);
  assert(dv.size() == a.size());
  for (int i = 0; i < dv.size(); ++i)
    dv[i] = (a[i] + kMod - dv[i]) % kMod;
  while (!dv.empty() && dv.back() == 0) dv.pop_back();
  return dv;
Poly Integral(const Poly &f) {
  int n = f.size();
  VI res(n + 1);
  for (int i = 0; i < n; ++i)
    res[i+1] = 1LL * f[i] * fpow(i + 1, kMod - 2)%kMod;
  return res;
Poly Evaluate(const Poly &f, const VI &x) {
  if (x.empty()) return Poly();
  int n = x.size();
  vector<Poly> up(n * 2);
  for (int i = 0; i < n; ++i) up[i+n] = {kMod-x[i], 1};
  for (int i = n - 1; i > 0; --i)
    up[i] = Multiply(up[i * 2], up[i * 2 + 1]);
  vector<Poly> down(n * 2)
  down[1] = Modulo(f, up[1]);
  for (int i = 2; i < n * 2; ++i)</pre>
    down[i] = Modulo(down[i >> 1], up[i]);
  VI y(n);
  for (int i = 0; i < n; ++i) y[i] = down[i + n][0];
  return y;
Poly Interpolate(const VI &x, const VI &y) {
  int n = x.size();
  vector<Poly> up(n * 2);
  for (int i = 0; i < n; ++i) up[i+n] = {kMod-x[i], 1};
  for (int i = n - 1; i > 0; --i)
  up[i] = Multiply(up[i * 2], up[i * 2 + 1]);
  VI a = Evaluate(Derivative(up[1]), x);
  for (int i = 0; i < n; ++i)
    a[i] = 1LL * y[i] * fpow(a[i], kMod - 2) % kMod;
  vector<Poly> down(n * 2);
  for (int i = 0; i < n; ++i) down[i + n] = {a[i]};
```

```
for (int i = n - 1; i > 0; --i) {
  auto lhs = Multiply(down[i * 2], up[i * 2 + 1]);
    auto rhs = Multiply(down[i * 2 + 1], up[i * 2]);
    assert(lhs.size() == rhs.size());
    down[i].resize(lhs.size());
    for (int j = 0; j < lhs.size(); ++j)</pre>
      down[i][j] = (lhs[j] + rhs[j]) % kMod;
  return down[1]:
Poly Log(Poly f) {
  int n = f.size();
  if (n == 1) return {0};
  auto d = Derivative(f);
  f.resize(n - 1)
  d = Multiply(d, Inverse(f));
  d.resize(n - 1);
  return Integral(d);
Poly Exp(Poly f) {
  int n = f.size()
  Poly q(1, 1); f[0] += 1;
  for (int s = 1; s < n; s <<= 1) {
   if (f.size() < s + s) f.resize(s + s);</pre>
    Poly g(f.begin(), f.begin() + s + s);
    Poly h(q.begin(), q.end());
h.resize(s + s); h = Log(h);
    for (int i = 0; i < s + s; ++i)
      g[i] = (g[i] + kMod - h[i]) % kMod;
    g = Multiply(g, q);
    g.resize(s + s); q = g;
  assert(q.size() >= n);
  q.resize(n);
  return q;
Poly SquareRootImpl(Poly f) {
  if (f.empty()) return {0};
  int z = QuadraticResidue(f[0], kMod), n = f.size();
  constexpr int kInv2 = (kMod + 1) >> 1;
  if (z == -1) return {-1};
  VI q(1, z);
  for (int s = 1; s < n; s <<= 1) {
    if (f.size() < s + s) f.resize(s + s);</pre>
    VI fq(q.begin(), q.end());
    fq.resize(s + s);
    VI f2 = Multiply(fq, fq);
    f2.resize(s + s);
    for (int i = 0; i < s + s; ++i)
f2[i] = (f2[i] + kMod - f[i]) % kMod;
    f2 = Multiply(f2, Inverse(fq));
    f2.resize(s + s);
    for (int i = 0; i < s + s; ++i)
      fq[i] = (fq[i]+kMod - 1LL*f2[i]*kInv2%kMod)%kMod;
    q = fq;
  q.resize(n);
  return q;
Poly SquareRoot(Poly f) {
  int n = f.size(), m = 0;
  while (m < n \&\& f[m] == 0) m++;
  if (m == n) return VI(n);
  if (m & 1) return {-1}
  auto s = SquareRootImpl(VI(f.begin() + m, f.end()));
  if (s[0] == -1) return {-1};
  VI res(n);
  for (int i = 0; i < s.size(); ++i) res[i + m/2]=s[i];</pre>
  return res;
```

#### 5.17 FWT

```
/* xor convolution:

* x = (x0,x1) , y = (y0,y1)

* z = ( x0y0 + x1y1 , x0y1 + x1y0 )

* =>

* x' = ( x0+x1 , x0-x1 ) , y' = ( y0+y1 , y0-y1 )

* z' = ( ( x0+x1 )( y0+y1 ) , ( x0-x1 )( y0-y1 ) )

* z = (1/2) * z''

* or convolution:
```

```
* x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
 * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
const LL MOD = 1e9+7;
inline void fwt( LL \times [MAXN] , int N , bool inv=0 ) {
  for( int d = 1 ; d < N ; d <<= 1 ) {</pre>
    int d2 = d << 1
    for( int s = 0 ; s < N ; s += d2 )
       for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
  LL ta = x[ i ] , tb = x[ j ];</pre>
         x[ i ] = ta+tb;
x[ j ] = ta-tb;
         if( x[ i ] >= MOD ) x[ i ] -= MOD;
         if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
       }
  if( inv )
    for( int i = 0 ; i < N ; i++ ) {</pre>
      x[ i ] *= inv( N, MOD );
       x[ i ] %= MOD;
}
```

### 5.18 DiscreteLog

```
// Baby-step Giant-step Algorithm
11d BSGS(11d P, 11d B, 11d N) {
  // find B^L = N \mod P
  unordered_map<lld, int> R;
  1ld sq = (lld)sqrt(P);
  11d t = 1;
  for (int i = 0; i < sq; i++) {
    if (t == N) return i
    if (!R.count(t)) R[t] = i;
    t = (t * B) % P;
  11d f = inverse(t, P);
  for(int i=0;i<=sq+1;i++) {</pre>
    if (R.count(N))
      return i * sq + R[N];
    N = (N * f) % P;
  return -1;
```

## 5.19 Quadratic residue

```
struct Status{
  11 x,y;
11 w:
Status mult(const Status& a,const Status& b,ll mod){
  Status res:
  res.x=(a.x*b.x+a.y*b.y%mod*w)%mod;
  res.y=(a.x*b.y+a.y*b.x)%mod;
  return res;
inline Status qpow(Status _base, 11 _pow, 11 _mod) {
  Status res = \{1, 0\};
  while(_pow>0){
    if(_pow&1) res=mult(res,_base,_mod);
_base=mult(_base,_base,_mod);
    _pow>>=1;
  }
  return res;
inline 11 check(11 x,11 p){
  return qpow_mod(x,(p-1)>>1,p);
inline 11 get_root(11 n,11 p){
  if(p==2) return 1;
  if(check(n,p)==p-1) return -1;
  11 a;
  while(true){
    a=rand()%p;
    w=((a*a-n)%p+p)%p;
    if(check(w,p)==p-1) break;
  Status res = \{a, 1\}
  res=qpow(res,(p+1)>>1,p);
  return res.x;
```

```
5.20 De-Bruijn
```

| }

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
 if (t > n) {
   if (n % p == 0)
      for (int i = 1; i <= p; ++i)
        res[sz++] = aux[i];
 } else ·
    aux[t] = aux[t - p];
    db(t + 1, p, n, k);
    for (int i = aux[t - p] + 1; i < k; ++i) {
     aux[t] = i;
     db(t + 1, t, n, k);
   }
 }
int de_bruijn(int k, int n) {
 // return cyclic string of len k^n s.t. every string
  // of len n using k char appears as a substring.
 if (k == 1) {
   res[0] = 0;
    return 1;
 for (int i = 0; i < k * n; i++) aux[i] = 0;
 db(1, 1, n, k);
 return sz;
```

## 5.21 Simplex Construction

Standard form: maximize  $\sum_{1\leq i\leq n} c_i x_i$  such that for all  $1\leq j\leq m$ ,  $\sum_{1\leq i\leq n} A_{ji} x_i \leq b_j$  and  $x_i\geq 0$  for all  $1\leq i\leq n$ .

- 1. In case of minimization, let  $c_i^\prime = -c_i$
- 2.  $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j$
- 3.  $\sum_{1 < i < n} A_{ji} x_i = b_j$ 
  - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$
  - $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

#### 5.22 Simplex

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return vector<double>(n, -inf) if the solution doesn
    't exist
// return vector<double>(n, +inf) if the solution is
    unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d:
vector<int> p, q;
void pivot(int r, int s) {
  double inv = 1.0 / d[r][s];
  for (int i = 0; i < m + 2; ++i) {
    for (int j = 0; j < n + 2; ++j) {
  if (i != r && j != s)</pre>
        d[i][j] -= d[r][j] * d[i][s] * inv;
   }
  for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;</pre>
  for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
  d[r][s] = inv;
 swap(p[r], q[s]);
bool phase(int z) {
 int x = m + z:
  while (true) {
    int s = -1;
    for (int i = 0; i <= n; ++i) {
```

```
if (!z && q[i] == -1) continue;
        if (s == -1) | d[x][i] < d[x][s]) s = i;
     if (d[x][s] > -eps) return true;
     int r = -1;
     for (int i = 0; i < m; ++i) {
       if (d[i][s] < eps) continue;
if (r == -1 || \</pre>
          d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
     if (r == -1) return false;
     pivot(r, s);
  }
VD solve(const VVD &a, const VD &b, const VD &c) {
   m = b.size(), n = c.size();
   d = VVD(m + 2, VD(n + 2));
for (int i = 0; i < m; ++i) {
     for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
   p.resize(m), q.resize(n + 1);
   for (int i = 0; i < m; ++i)</pre>
   p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
   q[n] = -1, d[m + 1][n] = 1;
   int r = 0;
   for (int i = 1; i < m; ++i)</pre>
     if(d[i][n+1] < d[r][n+1]) r = i;
   if (d[r][n + 1] < -eps) {</pre>
     pivot(r, n);
     if (!phase(1) || d[m + 1][n + 1] < -eps)
     return VD(n, -inf);
for (int i = 0; i < m; ++i) if (p[i] == -1) {
  int s = min_element(d[i].begin(), d[i].end() - 1)</pre>
                  - d[i].begin();
       pivot(i, s);
   if (!phase(0)) return VD(n, inf);
   VD x(n);
   for (int i = 0; i < m; ++i)
     if (p[i] < n) x[p[i]] = d[i][n + 1];
   return x;
}}
```

# 6 Geometry

#### 6.1 Point Class

```
template<typename T>
struct Point{
  typedef long double llf;
  static constexpr llf EPS = 1e-8;
  T x, y;
  Point(T _=0, T _{_}=0): x(_), y(_{_}){}
  template<typename T2>
    Point(const Point<T2>& a): x(a.x), y(a.y){}
  inline llf theta() const {
    return atan2((11f)y, (11f)x);}
  inline llf dis() const -
    return hypot((llf)x, (llf)y);}
  inline 11f dis(const Point& o) const {
    return hypot((11f)(x-o.x), (11f)(y-o.y));}
  Point operator-(const Point& o) const {
    return Point(x-o.x, y-o.y);}
  Point operator+(const Point& o) const {
 return Point(x+o.x, y+o.y);}
Point operator*(const T& k) const {
    return Point(x*k, y*k);}
  Point operator/(const T& k) const {
    return Point(x/k, y/k);}
  Point operator-() const {return Point(-x, -y);}
  Point rot90() const {return Point(-y, x);}
  template<typename T2>
  bool in(const Circle<T2>& a) const {
    /* Add struct Circle at top */
    return a.o.dis(*this)+EPS <= a.r;
  bool equal(const Point& o, true_type) const {
    return fabs(x-o.x) < EPS and fabs(y-o.y) < EPS; }</pre>
  bool equal(const Point& o, false_type) const {
    return tie(x, y) == tie(o.x, o.y); }
```

```
bool operator==(const Point& o) const {
  return equal(o, is_floating_point<T>()); }
bool operator!=(const Point& o) const {
  return !(*this == o); }
bool operator<(const Point& o) const {</pre>
  return theta() < o.theta();</pre>
  // sort like what pairs did
  // if (is_floating_point<T>())
      return fabs(x-o.x)<EPS?y<o.y:x<o.x;
  // else return tie(x, y) < tie(o.x, o.y);</pre>
friend inline T cross(const Point&a,const Point&b){
  return a.x*b.y - b.x*a.y; }
friend inline T dot(const Point& a, const Point &b){
  return a.x*b.x + a.y*b.y; }
friend ostream&operator<<(ostream&ss,const Point&o){</pre>
  ss<<"("<<o.x<<", "<<o.y<<")"; return ss; }
```

## 6.2 Circle Class

```
template<typename T>
struct Circle{
    static constexpr llf EPS = 1e-8;
    Point<T> o; T r;
    vector<Point<llf>> operator&(const Circle& aa)const{
        //https://www.cnblogs.com/wangzming/p/8338142.html
        llf d=o.dis(aa.o);
        if(d>r+aa.r+EPS || d<fabs(r-aa.r)-EPS) return {};
        llf dt = (r*r - aa.r*aa.r)/d, d1 = (d+dt)/2;
        Point<llf>> dir = (aa.o-o); dir /= d;
        Point<llf>> pcrs = dir*d1 + o;
        dt=sqrt(max(0.0L, r*r - d1*d1)), dir=dir.rot90();
        return {pcrs + dir*dt, pcrs - dir*dt};
    }
};
```

## 6.3 Segment Class

```
const long double EPS = 1e-8;
template<typename T>
struct Segment{
  // p1.x < p2.x
 Line<T> base;
 Point<T> p1, p2;
 Segment(): base(Line<T>()), p1(Point<T>()), p2(Point<</pre>
      T>()){
   assert(on_line(p1, base) and on_line(p2, base));
 Segment(Line<T> _, Point<T> __, Point<T> __): base(_
      ), p1(__), p2(___){
    assert(on_line(p1, base) and on_line(p2, base));
 template<typename T2>
    Segment(const Segment<T2>& _): base(_.base), p1(_.
        p1), p2(_.p2) {}
  typedef Point<long double> Pt;
  friend bool on_segment(const Point<T>& p, const
      Segment& 1) {
    if(on_line(p, 1.base))
      return (1.p1.x-p.x)*(p.x-1.p2.x)>=0 and (1.p1.y-p.x)
          .y)*(p.y-1.p2.y)>=0;
    return false;
 friend bool have_inter(const Segment& a, const
      Segment& b) {
    if(is_parallel(a.base, b.base)){
      return on_segment(a.p1, b) or on_segment(a.p2, b)
           or on_segment(b.p1, a) or on_segment(b.p2, a
          ):
    Pt inter = get_inter(a.base, b.base);
    return on_segment(inter, a) and on_segment(inter, b
  friend inline Pt get_inter(const Segment& a, const
      Segment& b) {
    if(!have_inter(a, b)){
      return NOT_EXIST;
    }else if(is_parallel(a.base, b.base)){
     if(a.p1 == b.p1){
```

```
if(on_segment(a.p2, b) or on_segment(b.p2, a))
            return INF_P;
        else return a.p1;
      else if(a.p1 == b.p2){
        if(on_segment(a.p2, b) or on_segment(b.p1, a))
            return INF_P;
        else return a.p1;
      else if(a.p2 == b.p1){
        if(on_segment(a.p1, b) or on_segment(b.p2, a))
            return INF_P
        else return a.p2;
      }else if(a.p2 == b.p2){
        if(on_segment(a.p1, b) or on_segment(b.p1, a))
            return INF_P;
        else return a.p2;
      return INF_P;
    return get_inter(a.base, b.base);
  friend ostream& operator<<(ostream& ss, const Segment</pre>
      & o){
    ss<<o.base<<", "<<o.p1<<" ~ "<<o.p2;
    return ss;
template<typename T>
inline Segment<T> get_segment(const Point<T>& a, const
    Point<T>& b){
  return Segment<T>(get_line(a, b), a, b);
```

#### 6.4 Line Class

```
const Point<long double> INF_P(-1e20, 1e20);
const Point<long double> NOT_EXIST(1e20, 1e-20);
template<typename T>
struct Line{
 static constexpr long double EPS = 1e-8;
  // ax+by+c = 0
 T a, b, c;
 Line(T _=0, T __=1, T __=0): a(_), b(_), c(_){
    assert(fabs(a)>EPS or fabs(b)>EPS);}
  template<typename T2>
    Line(const Line<T2>& x): a(x.a), b(x.b), c(x.c){}
  typedef Point<long double> Pt;
  bool equal(const Line& o, true_type) const {
    return fabs(a-o.a)<EPS &&
    fabs(b-o.b)<EPS && fabs(c-o.b)<EPS;}
  bool equal(const Line& o, false_type) const {
    return a==o.a and b==o.b and c==o.c;}
  bool operator==(const Line& o) const {
    return equal(o, is_floating_point<T>());}
  bool operator!=(const Line& o) const {
    return !(*this == o);}
  friend inline bool on_line__(const Point<T>& p, const
       Line& 1, true_type){
    return fabs(1.a*p.x + 1.b*p.y + 1.c) < EPS;
  friend inline bool on_line__(const Point<T>& p, const
       Line& 1, false_type){
    return 1.a*p.x + 1.b*p.y + 1.c == 0;
  friend inline bool on_line(const Point<T>&p, const
      Line& 1){
    return on_line__(p, 1, is_floating_point<T>());
  friend inline bool is_parallel__(const Line& x, const
       Line& y, true_type){
    return fabs(x.a*y.b - x.b*y.a) < EPS;</pre>
  friend inline bool is_parallel__(const Line& x, const
       Line& y, false_type){
    return x.a*y.b == x.b*y.a;
  friend inline bool is_parallel(const Line& x, const
      Line& y){
    return is_parallel__(x, y, is_floating_point<T>());
  friend inline Pt get_inter(const Line& x, const Line&
       y){
```

```
typedef long double llf;
if(x==y) return INF_P;
if(is_parallel(x, y)) return NOT_EXIST;
llf delta = x.a*y.b - x.b*y.a;
llf delta_x = x.b*y.c - x.c*y.b;
llf delta_y = x.c*y.a - x.a*y.c;
return Pt(delta_x / delta, delta_y / delta);
}
friend ostream&operator<<(ostream&ss, const Line&o){
    s<<0.a<<"x+"<<0.b<<"y+"<<0.c<<"=0";
    return ss;
}
};
template<typename T>
inline Line<T> get_line(const Point<T>& a, const Point<
    T>& b){
    return Line<T>(a.y-b.y, b.x-a.x, (b.y-a.y)*a.x-(b.x-a.x)*a.y);
}
```

## 6.5 Triangle Circumcentre

### 6.6 2D Convex Hull

```
template<typename T>
class ConvexHull_2D{
private:
 typedef Point<T> PT;
 vector<PT> d;
 struct myhash{
   uint64_t operator()(const PT& a) const {
      uint64_t xx=0, yy=0;
      memcpy(&xx, &a.x, sizeof(a.x));
      memcpy(&yy, &a.y, sizeof(a.y));
      uint64_t ret = xx*17+yy*31;
      ret = (ret ^ (ret >> 16))*0x9E3779B1;
      ret = (ret ^ (ret >> 13))*0xC2B2AE35;
      ret = ret ^ xx;
      return (ret ^ (ret << 3)) * yy;</pre>
 unordered_set<PT, myhash> in_hull;
public:
 void init(){in_hull.clear();d.clear();}
  void insert(const PT& x){d.PB(x);}
 void solve(){
    sort(ALL(d), [](const PT& a, const PT& b){
      return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
    vector<PT> s(SZ(d)<<1); int o = 0;
    for(auto p: d) {
      while(o \ge 2 \& cross(p-s[o-2], s[o-1]-s[o-2]) <= 0)
        0--
      s[o++] = p;
    for(int i=SZ(d)-2, t = o+1;i>=0;i--){
      while(o = t\&cross(d[i] - s[o-2], s[o-1] - s[o-2]) <= 0)
      s[o++] = d[i];
    s.resize(o-1); swap(s, d);
    for(auto i: s) in_hull.insert(i);
 vector<PT> get(){return d;}
 bool in_it(const PT& x){
    return in_hull.find(x)!=in_hull.end();}
```

## 6.7 2D Farthest Pair

```
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {
  while(abs(cross(stk[i+1]-stk[i],
        stk[(pos+1)%n]-stk[i])) >
        abs(cross(stk[i+1]-stk[i],
        stk[pos]-stk[i]))) pos = (pos+1)%n;
ans = max({ans, dis(stk[i], stk[pos]),
        dis(stk[i+1], stk[pos]));
}
```

#### 6.8 2D Closest Pair

```
struct Pt{
  11f x, y, d;
} arr[N];
inline llf dis(Pt a, Pt b){
  return hypot(a.x-b.x, a.y-b.y);
11f solve(){
  int cur = rand() % n;
  for(int i=0;i<n;i++) arr[i].d = dis(arr[cur], arr[i])</pre>
  sort(arr, arr+n, [](Pt a, Pt b){return a.d < b.d;});</pre>
  11f ans = 1e50;
  for(int i=0;i<n;i++){</pre>
    for(int j=i+1;j<n;j++){</pre>
      if(arr[j].d - arr[i].d > ans) break;
      ans = min(ans, dis(arr[i], arr[j]));
  return ans;
}
```

## 6.9 Simulated Annealing

```
Ilf anneal() {
  mt19937 rnd_engine( seed );
  uniform_real_distribution< llf > rnd( 0, 1 );
  const llf dT = 0.001;
  // Argument p
  llf S_cur = calc( p ), S_best = S_cur;
  for ( llf T = 2000 ; T > EPS ; T -= dT ) {
      // Modify p to p_prime
      const llf S_prime = calc( p_prime );
      const llf delta_c = S_prime - S_cur;
      llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
      if ( rnd( rnd_engine ) <= prob )
            S_cur = S_prime, p = p_prime;
      if ( S_prime < S_best )
            S_best = S_prime, p_best = p_prime;
    }
    return S_best;
}</pre>
```

## 6.10 Half Plane Intersection

```
inline int dcmp ( double x ) {
  if( fabs( x ) < eps ) return 0;</pre>
  return x > 0 ? 1 : -1;
struct Line {
  Point st, ed;
  double ang;
  Line(Point _s=Point(), Point _e=Point()):
   st(_s),ed(_e),ang(atan2(_e.y-_s.y,_e.x-_s.x)){}
  inline bool operator< ( const Line& rhs ) const {</pre>
    if(dcmp(ang - rhs.ang) != 0) return ang < rhs.ang;</pre>
    return dcmp( cross( st, ed, rhs.st ) ) < 0;</pre>
// cross(pt, line.ed-line.st)>=0 <-> pt in half plane
vector< Line > lns;
deque< Line > que;
deque< Point > pt;
double HPI() {
  sort( lns.begin(), lns.end() );
```

```
que.clear(); pt.clear();
que.push_back( lns[ 0 ] );
for ( int i = 1 ; i < (int)lns.size() ; i ++ ) {</pre>
  if(!dcmp(lns[i].ang - lns[i-1].ang)) continue;
  while ( pt.size() > 0 &&
   dcmp(cross(lns[i].st,lns[i].ed,pt.back()))<0){</pre>
    pt.pop_back();que.pop_back();
  while ( pt.size() > 0 &&
   dcmp(cross(lns[i].st,lns[i].ed,pt.front()))<0){</pre>
    pt.pop_front(); que.pop_front();
  pt.push_back(get_point( que.back(), lns[ i ] ));
  que.push_back( lns[ i ] );
while ( pt.size() > 0 &&
 dcmp(cross(que[0].st, que[0].ed, pt.back()))<0){</pre>
  que.pop_back();
  pt.pop_back();
while ( pt.size() > 0 &&
 dcmp(cross(que.back().st,que.back().ed,pt[0]))<0){</pre>
  que.pop_front();
  pt.pop_front();
pt.push_back(get_point(que.front(), que.back()));
vector< Point > conv;
for ( int i = 0 ; i < (int)pt.size() ; i ++ )</pre>
  conv.push_back( pt[ i ] );
double ret = 0;
for ( int i = 1 ; i + 1 < (int)conv.size() ; i ++ )</pre>
  ret += abs(cross(conv[0], conv[i], conv[i + 1]));
return ret / 2.0;
```

## 6.11 Ternary Search on Integer

```
int TernarySearch(int 1, int r) {
    // max value @ (1, r]
    while (r - 1 > 1){
        int m = (1 + r)>>1;
        if (f(m) > f(m + 1)) r = m;
        else 1 = m;
    }
    return 1+1;
}
```

## 6.12 Minimum Covering Circle

```
template<typename T>
Circle<llf> MinCircleCover(const vector<PT>& pts){
  random_shuffle(ALL(pts));
  Circle<llf> c = {pts[0], 0};
  for(int i=0;i<SZ(pts);i++){
    if(pts[i].in(c)) continue;
    c = {pts[i], 0};
  for(int j=0;j<i;j++){
      if(pts[j].in(c)) continue;
      c.o = (pts[i] + pts[j]) / 2;
      c.r = pts[i].dis(c.o);
    for(int k=0;k<j;k++){
      if(pts[k].in(c)) continue;
      c = get_circum(pts[i], pts[j], pts[k]);
    }
  }
}
return c;
}</pre>
```

## 6.13 KDTree (Nearest Point)

```
const int MXN = 100005;
struct KDTree {
    struct Node {
        int x,y,x1,y1,x2,y2;
        int id,f;
        Node *L, *R;
    } tree[MXN], *root;
    int n;
    LL dis2(int x1, int y1, int x2, int y2) {
```

```
LL dx = x1-x2, dy = y1-y2;
     return dx*dx+dy*dy;
  static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
  static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
  void init(vector<pair<int,int>> ip) {
    n = ip.size();
     for (int i=0; i<n; i++) {
       tree[i].id = i;
       tree[i].x = ip[i].first;
       tree[i].y = ip[i].second;
     root = build_tree(0, n-1, 0);
  Node* build_tree(int L, int R, int d) {
     if (L>R) return nullptr;
     int M = (L+R)/2; tree[M].f = d%2;
     nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
     tree[M].x1 = tree[M].x2 = tree[M].x;
     tree[M].y1 = tree[M].y2 = tree[M].y;
     tree[M].L = build_tree(L, M-1, d+1);
     if (tree[M].L) {
       tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
       tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
       tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
     tree[M].R = build_tree(M+1, R, d+1);
     if (tree[M].R) {
       tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
     return tree+M;
  int touch(Node* r, int x, int y, LL d2){
    LL dis = sqrt(d2)+1;
     if (x<r->x1-dis || x>r->x2+dis ||
         y<r->y1-dis || y>r->y2+dis)
       return 0;
    return 1;
  void nearest(Node* r,int x,int y,int &mID,LL &md2) {
     if (!r || !touch(r, x, y, md2)) return;
     LL d2 = dis2(r->x, r->y, x, y);
    if (d2 < md2 || (d2 == md2 && mID < r->id)) {
      mID = r -> id;
      md2 = d2;
     // search order depends on split dim
    if ((r->f == 0 \&\& x < r->x) ||
         (r->f == 1 \&\& y < r->y))
       nearest(r->L, x, y, mID, md2);
       nearest(r->R, x, y, mID, md2);
     } else {
       nearest(r->R, x, y, mID, md2);
       nearest(r->L, x, y, mID, md2);
  int query(int x, int y) {
    int id = 1029384756;
     LL d2 = 102938475612345678LL;
    nearest(root, x, y, id, d2);
    return id;
} tree;
```

# 7 Stringology

#### 7.1 Hash

```
class Hash{
private:
    const int p = 127, q = 1051762951;
    int sz, prefix[N], power[N];
    int add(int x, int y){return x+y>=q?x+y-q:x+y;}
    int sub(int x, int y){return x-y<0?x-y+q:x-y;}
    int mul(int x, int y){return 1LL*x*y%q;}
public:
    void init(const string &x){</pre>
```

```
sz = x.size();prefix[0]=0;power[0]=1;
for(int i=1;i<=sz;i++)
    prefix[i]=add(mul(prefix[i-1], p), x[i-1]);
for(int i=1;i<=sz;i++)power[i]=mul(power[i-1], p);
}
int query(int l, int r){
    // 1-base (l, r]
    return sub(prefix[r], mul(prefix[l], power[r-l]));
}
};</pre>
```

# 7.2 Suffix Array

```
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
int x[maxn], p[maxn], q[maxn * 2];
// sa[i]: sa[i]-th suffix is the \
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix \
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *sa, int *c, int n, int z) {
  memset(sa, 0, sizeof(int) * n);
  memcpy(x, c, sizeof(int) * z);
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
 memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
if (sa[i] && !t[sa[i] - 1])
      sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  memcpy(x, c, sizeof(int) * z);
for (int i = n - 1; i >= 0; --i)
    if (sa[i] && t[sa[i] - 1])
      sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q,
 bool *t, int *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
  memset(c, 0, sizeof(int) * z);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
  for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
    return;
  for (int i = n - 2; i >= 0; --i)
t[i] = (s[i]==s[i + 1] ? t[i + 1] : s[i]<s[i + 1]);
  pre(sa, c, n, z);
  for (int i = 1; i <= n - 1; ++i)
    if (t[i] && !t[i - 1])
      sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) {
    if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
    bool neq = last < 0 || \</pre>
     memcmp(s + sa[i], s + last,
(p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
    ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
  pre(sa, c, n, z);
  for (int i = nn - 1; i >= 0; --i)
    sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
void build(const string &s) {
  for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
   _s[(int)s.size()] = 0; // s shouldn't contain 0
  sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256)
  for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];
  for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
  int ind = 0; hi[0] = 0;
  for (int i = 0; i < (int)s.size(); ++i) {</pre>
    if (!rev[i]) {
      ind = 0;
      continue;
    while (i + ind < (int)s.size() && \</pre>
     s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
```

```
|}}
```

hi[rev[i]] = ind ? ind-- : 0;

```
7.3
       Aho-Corasick Algorithm
class AhoCorasick{
  private:
    static constexpr int Z = 26;
    struct node{
      node *nxt[ Z ], *fail;
      vector< int > data;
      node(): fail( nullptr ) {
        memset( nxt, 0, sizeof( nxt ) );
         data.clear();
    } *rt;
    inline int Idx( char c ) { return c - 'a'; }
  public:
    void init() { rt = new node(); }
    void add( const string& s, int d ) {
      node* cur = rt;
      for ( auto c : s ) {
        if ( not cur->nxt[ Idx( c ) ] )
          cur->nxt[ Idx( c ) ] = new node();
        cur = cur->nxt[ Idx( c ) ];
      cur->data.push_back( d );
    void compile() {
  vector< node* > bfs;
      size_t ptr = 0;
      for ( int i = 0 ; i < Z ; ++ i ) {
  if ( not rt->nxt[ i ] ) {
          // uncomment 2 lines to make it DFA
           // rt->nxt[i] = rt;
          continue;
         rt->nxt[ i ]->fail = rt;
        bfs.push_back( rt->nxt[ i ] );
      while ( ptr < bfs.size() ) {</pre>
        node* u = bfs[ ptr ++ ];
for ( int i = 0 ; i < Z ; ++ i ) {</pre>
           if ( not u->nxt[ i ] ) {
             // u->nxt[i] = u->fail->nxt[i];
             continue;
          node* u_f = u->fail;
           while ( u_f ) {
             if ( not u_f->nxt[ i ] ) {
               u_f = u_f->fail; continue;
             u->nxt[ i ]->fail = u_f->nxt[ i ];
           if ( not u_f ) u->nxt[ i ]->fail = rt;
          bfs.push_back( u->nxt[ i ] );
      }
    void match( const string& s, vector< int >& ret ) {
      node* u = rt;
      for ( auto c : s ) {
        while ( u != rt and not u->nxt[ Idx( c ) ] )
          u = u->fail;
         u = u - nxt[Idx(c)];
        if ( not u ) u = rt;
         node* tmp = u;
        while ( tmp != rt ) {
           for ( auto d : tmp->data )
             ret.push_back( d );
           tmp = tmp->fail;
      }
} ac;
```

#### 7.4 Suffix Automaton

for (int i = 0, k = 0; i < (int)s.size(); ++i)</pre>

k = f[k - 1]

if (s[i] == t[k]) ++k;

 $while(k > 0 \&\& (k==(int)t.size() \mid | s[i]!=t[k]))$ 

if (k == (int)t.size()) r.push\_back(i-t.size()+1);

```
struct Node{
  Node *green, *edge[26];
                                                                 return res;
  int max_len;
  Node(const int _max_len)
    : green(NULL), max_len(_max_len){
                                                               7.6 Z value
    memset(edge,0,sizeof(edge));
                                                               char s[MAXN];
} *ROOT, *LAST;
                                                               int len, z[MAXN];
void Extend(const int c) {
                                                               void Z_value() {
  Node *cursor = LAST;
                                                                 int i, j, left, right;
  LAST = new Node((LAST->max_len) + 1);
                                                                 left=right=0; z[0]=len;
  for(;cursor&&!cursor->edge[c]; cursor=cursor->green)
                                                                 for(i=1;i<len;i++) {</pre>
    cursor->edge[c] = LAST;
                                                                   j=max(min(z[i-left], right-i),0);
  if (!cursor)
                                                                    for(;i+j<len&&s[i+j]==s[j];j++);
    LAST->green = ROOT;
                                                                   z[i]=j;
  else {
                                                                   if(i+z[i]>right) {
    Node *potential_green = cursor->edge[c];
                                                                     right=i+z[i];
    if((potential_green->max_len)==(cursor->max_len+1))
                                                                     left=i;
      LAST->green = potential_green;
    else {
//assert(potential_green->max_len>(cursor->max_len+1));
      Node *wish = new Node((cursor->max_len) + 1);
      for(;cursor && cursor->edge[c]==potential_green;
            cursor = cursor->green)
                                                               7.7 Manacher
        cursor->edge[c] = wish;
      for (int i = 0; i < 26; i++)
  wish->edge[i] = potential_green->edge[i];
                                                               int z[maxn];
                                                               int manacher(const string& s) {
      wish->green = potential_green->green;
                                                                 string t =
      potential_green->green = wish;
                                                                 for(char c:s)) t += c, t += '.';
      LAST->green = wish;
                                                                 int 1 = 0, r = 0, ans = 0;
    }
                                                                 for (int i = 1; i < t.length(); ++i) {</pre>
  }
                                                                   z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
while (i - z[i] >= 0 && i + z[i] < t.length()) {
char S[10000001], A[10000001];
                                                                     if(t[i - z[i]] == t[i + z[i]]) ++z[i];
int N;
                                                                      else break;
int main(){
  scanf("%d%s", &N, S);
                                                                   if (i + z[i] > r) r = i + z[i], l = i;
  ROOT = LAST = new Node(0);
  for (int i = 0; S[i]; i++)
                                                                 for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
    Extend(S[i] - 'a');
                                                                 return ans;
  while (N--){
    scanf("%s", A);
    Node *cursor = ROOT;
    bool ans = true;
                                                                      Lexico Smallest Rotation
    for (int i = 0; A[i]; i++){
      cursor = cursor->edge[A[i] - 'a'];
                                                               string mcp(string s){
      if (!cursor) {
                                                                 int n = s.length();
        ans = false;
                                                                 s += s;
        break;
                                                                 int i=0, j=1;
      }
                                                                 while (i<n && j<n){</pre>
                                                                   int k = 0;
    puts(ans ? "Yes" : "No");
                                                                   while (k < n \&\& s[i+k] == s[j+k]) k++;
                                                                   if (s[i+k] <= s[j+k]) j += k+1;
  return 0;
                                                                   else i += k+1;
                                                                   if (i == j) j++;
                                                                 int ans = i < n ? i : j;</pre>
7.5 KMP
                                                                 return s.substr(ans. n):
                                                               }
vector<int> kmp(const string &s) {
  vector<int> f(s.size(), 0);
                                                               7.9
                                                                      BWT
  /* f[i] = length of the longest prefix
     (excluding s[0:i]) such that it coincides
     with the suffix of s[0:i] of the same length */
                                                               struct BurrowsWheeler{
  /*i+1-f[i] is the length of the
                                                               #define SIGMA 26
     smallest recurring period of s[0:i] */
                                                               #define BASE 'a'
  int k = 0:
                                                                 vector<int> v[ SIGMA ];
  for (int i = 1; i < (int)s.size(); ++i) {
  while (k > 0 && s[i] != s[k]) k = f[k - 1];
                                                                 void BWT(char* ori, char* res){
  // make ori -> ori + ori
     if (s[i] == s[k]) ++k;
                                                                   // then build suffix array
    f[i] = k;
                                                                 void iBWT(char* ori, char* res){
                                                                   for( int i = 0 ; i < SIGMA ; i ++ )</pre>
  return f;
                                                                     v[ i ].clear();
vector<int> search(const string &s, const string &t) {
                                                                    int len = strlen( ori );
                                                                   for( int i = 0 ; i < len ; i ++ )
  v[ ori[i] - BASE ].push_back( i );</pre>
  // return 0-indexed occurrence of t in s
  vector < int > f = kmp(t), r;
```

vector<int> a;

for( auto j : v[ i ] ){

a.push\_back( j ); ori[ ptr ++ ] = BASE + i;

for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>

```
}
for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
    res[ i ] = ori[ a[ ptr ] ];
    ptr = a[ ptr ];
}
res[ len ] = 0;
}
} bwt;
</pre>
```

#### 7.10 Palindromic Tree

```
struct palindromic_tree{
  struct node{
    int next[26],f,len;
    int cnt,num,st,ed;
    node(int l=0):f(0),len(1),cnt(0),num(0) {
      memset(next, 0, sizeof(next)); }
  }:
  vector<node> st;
  vector<char> s;
  int last,n;
  void init(){
    st.clear();s.clear();last=1; n=0;
    st.push_back(0);st.push_back(-1);
    st[0].f=1;s.push_back(-1); }
  int getFail(int x){
    while(s[n-st[x].len-1]!=s[n])x=st[x].f;
    return x;}
  void add(int c){
    s.push_back(c-='a'); ++n;
    int cur=getFail(last);
    if(!st[cur].next[c]){
      int now=st.size();
      st.push_back(st[cur].len+2);
      st[now].f=st[getFail(st[cur].f)].next[c];
      st[cur].next[c]=now;
      st[now].num=st[st[now].f].num+1;
    last=st[cur].next[c];
    ++st[last].cnt;}
  int size(){ return st.size()-2;}
} pt;
int main() {
  string s; cin >> s; pt.init();
for (int i=0; i<SZ(s); i++) {</pre>
    int prvsz = pt.size(); pt.add(s[i]);
    if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
      // pal @ [1,r]: s.substr(1, r-1+1)
    }
  return 0;
```

## 8 Misc

## 8.1 Theorems

#### 8.1.1 Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G , where  $L_{ii}=d(i)$  ,  $L_{ij}=-c$  where c is the number of edge (i,j) in G .

- The number of undirected spanning in G is  $|\mathsf{det}(\tilde{L}_{11})|$  .
- The number of directed spanning tree rooted at r in G is  $|{\rm det}(\tilde{L}_{rr})|.$

#### 8.1.2 Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$   $(x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

#### 8.1.3 Cayley's Formula

- Given a degree sequence  $d_1,d_2,\ldots,d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .

#### 8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all  $1 \le k \le n$ .

#### 8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

#### 8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let  $N_G(W)$  denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff  $\forall W\subseteq X, |W|\leq |N_G(W)|$ 

### 8.1.7 Euler's planar graph formula

V - E + F = C + 1,  $E \le 3V - 6$ (?)

#### 8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have  $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$ 

## 8.2 MaximumEmptyRect

```
int max_empty_rect(int n, int m, bool blocked[N][N]) {
  static int mxu[2][N], me=0, he=1, ans=0;
  for (int i=0;i<m;i++) mxu[he][i]=0;</pre>
  for (int i=0;i<n;i++) {</pre>
    stack<PII, vector<PII>> stk;
     for (int j=0;j<m;++j) {</pre>
       if (blocked[i][j]) mxu[me][j]=0;
       else mxu[me][j]=mxu[he][j]+1;
       int la = j;
       while (!stk.empty()&&stk.top().FF>mxu[me][j]) {
         int x1 = i - stk.top().FF, x2 = i;
int y1 = stk.top().SS, y2 = j;
la = stk.top().SS; stk.pop();
         ans=max(ans,(x2-x1)*(y2-y1));
       if (stk.empty()||stk.top().FF<mxu[me][j])</pre>
         stk.push({mxu[me][j],la});
    while (!stk.empty()) {
       int x1 = i - stk.top().FF, x2 = i;
       int y1 = stk.top().SS-1, y2 = m-1;
       stk.pop(); ans=max(ans,(x2-x1)*(y2-y1));
    swap(me,he);
  return ans;
}
```

## 8.3 DP-opt Condition

## 8.3.1 totally monotone (concave/convex)

```
\begin{array}{c} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

## 8.3.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

## 8.4 Convex 1D/1D DP

```
struct segment {
   int i, 1, r;
   segment() {}
   segment(int a, int b, int c): i(a), l(b), r(c) {}
};
inline lld f(int l, int r){return dp[l] + w(l+1, r);}
void solve() {
   dp[0] = 0;
   deque<segment> dq; dq.push_back(segment(0, 1, n));
   for (int i = 1; i <= n; ++i) {
      dp[i] = f(dq.front().i, i);
}</pre>
```

```
while(dq.size()&&dq.front().r<i+1) dq.pop_front();
    dq.front().l = i + 1;
    segment seg = segment(i, i + 1, n);
    while (dq.size() &&
        f(i, dq.back().l)<f(dq.back().i, dq.back().l)
            dq.pop_back();
    if (dq.size()) {
        int d = 1 << 20, c = dq.back().l;
        while (d >>= 1) if (c + d <= dq.back().r)
            if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
        dq.back().r = c; seg.l = c + 1;
    }
    if (seg.l <= n) dq.push_back(seg);
}</pre>
```

## 8.5 ConvexHull Optimization

```
inline 1ld DivCeil(1ld n, 1ld d) { // ceil(n/d)
  return n / d + (((n < 0) != (d > 0)) && (n % d));
struct Line {
  static bool flag;
  11d a, b, 1, r; // y=ax+b in [1, r)
  11d operator()(11d x) const { return a * x + b; }
  bool operator<(const Line& i) const {</pre>
    return flag ? tie(a, b) < tie(i.a, i.b) : 1 < i.l;</pre>
  11d operator&(const Line& i) const {
    return DivCeil(b - i.b, i.a - a);
bool Line::flag = true;
class ConvexHullMax {
  set<Line> L;
 public:
  ConvexHullMax() { Line::flag = true; }
  void InsertLine(lld a, lld b) { // add y = ax + b
    Line now = \{a, b, -INF, INF\};
    if (L.empty())
      L.insert(now);
      return:
    Line::flag = true;
    auto it = L.lower_bound(now);
    auto prv = it == L.begin() ? it : prev(it);
    if (it != L.end() && ((it != L.begin() &&
      (*it)(it->1) >= now(it->1) &&
      (*prv)(prv->r-1) >= now(prv->r-1)) ||
      (it == L.begin() && it->a == now.a))) return;
    if (it != L.begin()) -
      while (prv != L.begin() &&
         (*prv)(prv->1) <= now(prv->1))
          prv = --L.erase(prv);
      if (prv == L.begin() && now.a == prv->a)
        L.erase(prv);
    if (it != L.end())
      while (it != --L.end() &&
         (*it)(it->r) <= now(it->r))
          it = L.erase(it);
    if (it != L.begin()) {
      prv = prev(it);
      const_cast<Line*>(&*prv)->r=now.l=((*prv)&now);
    if (it != L.end())
      const_cast<Line*>(&*it)->l=now.r=((*it)&now);
    L.insert(it, now);
  11d Query(11d a) const { // query max at x=a
    if (L.empty()) return -INF;
    Line::flag = false;
    auto it = --L.upper_bound(\{0, 0, a, 0\});
    return (*it)(a);
};
```

# 8.6 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
```

```
int s = 0;
for (int i = 2; i <= n; i++)
    s = (s + m) % i;
return s;
}
// died at kth
int kth(int n, int m, int k){
    if (m == 1) return n-1;
    for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
    return k;
}
```

# 8.7 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u)
  dfn[u]=low[u]=++dfs_idx;
  for(int i=0;i<(int)init_g[u].size();i++){</pre>
    int v=init_g[u][i];
    if(v==par[u]) continue;
    if(!dfn[v]){
      par[v]=u;
      tarian(v):
      low[u]=min(low[u],low[v]);
      if(dfn[u]<low[v]){</pre>
        g[u].push_back(v)
        g[v].push_back(u);
    }else{
      low[u]=min(low[u],dfn[v]);
      if(dfn[v]<dfn[u]){</pre>
        int temp_v=u;
        bcc_id++;
        while(temp_v!=v){
          g[bcc_id+n].push_back(temp_v);
          g[temp_v].push_back(bcc_id+n);
          temp_v=par[temp_v];
        g[bcc_id+n].push_back(v);
        g[v].push_back(bcc_id+n);
        reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
    }
  }
int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
void dfs(int u,int fa){
  if(u<=n){
    for(int i=0;i<(int)g[u].size();i++){</pre>
      int v=g[u][i];
      if(v==fa) continue;
      dfs(v,u);
      memset(tp,0x8f,sizeof tp);
      if(v<=n){
        tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
        tp[1]=max
          dp[u][0]+dp[v][0]+1
          dp[u][1]+max(dp[v][0],dp[v][1])
      }else{
        tp[0]=dp[u][0]+dp[v][0];
        tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
      dp[u][0]=tp[0],dp[u][1]=tp[1];
  }else{
    for(int i=0;i<(int)g[u].size();i++){</pre>
      int v=g[u][i];
      if(v==fa) continue;
      dfs(v,u);
    min_dp[0][0]=0;
    min_dp[1][1]=1;
    min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
    for(int i=0;i<(int)g[u].size();i++){</pre>
      int v=g[u][i];
      if(v==fa) continue;
      memset(tmp,0x8f,sizeof tmp);
      tmp[0][0]=max(
        \min_{dp[0][0]+\max(dp[v][0],dp[v][1])}
```

```
National Taiwan University - kiseki
         min_dp[0][1]+dp[v][0]
                                                                           ansd = d:
       tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
      tmp[1][0]=max(
         min_dp[1][0]+max(dp[v][0],dp[v][1]),
         min_dp[1][1]+dp[v][0]
       tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
      memcpy(min_dp,tmp,sizeof tmp);
    dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
    dp[u][0]=min_dp[0][0];
  }
int main(){
  int m,a,b;
  scanf("%d%d",&n,&m);
for(int i=0;i<m;i++){</pre>
                                                                       }
                                                                    } sol;
    scanf("%d%d",&a,&b);
    init_g[a].push_back(b);
    init_g[b].push_back(a);
 par[1]=-1;
  tarjan(1);
 dfs(1,-1);
  printf("%d\n", max(dp[1][0], dp[1][1]));
  return 0;
8.8
      DLX
struct DLX {
                                                                      }
  const static int maxn=210;
  const static int maxm=210;
  const static int maxnode=210*210;
  int n, m, size, row[maxnode], col[maxnode];
int U[maxnode], D[maxnode], L[maxnode], R[maxnode];
  int H[maxn], S[maxm], ansd, ans[maxn];
  void init(int _n, int _m) {
    n = _n, m = _m;
    for(int i = 0; i <= m; ++i) {</pre>
      S[i] = 0;
      U[i] = D[i] = i;
      L[i] = i-1, R[i] = i+1;
    R[L[0] = size = m] = 0;
    for(int i = 1; i <= n; ++i) H[i] = -1;
                                                                    8.10
  void Link(int r, int c) {
    ++S[col[++size] = c];
    row[size] = r; D[size] = D[c];
U[D[c]] = size; U[size] = c; D[c] = size;
if(H[r] < 0) H[r] = L[size] = R[size] = size;</pre>
    else {
      R[size] = R[H[r]];
      L[R[H[r]]] = size;
      L[size] = H[r];
      R[H[r]] = size;
    }
```

void remove(int c) {

D[U[j]] = D[j];

L[R[c]] = c; R[L[c]] = c;

--S[col[j]];

D[U[j]] = j++S[col[j]];

void dance(int d) { if(d>=ansd) return; **if**(R[0] == 0) {

} }

void resume(int c) {

L[R[c]] = L[c]; R[L[c]] = R[c];

for(int i = D[c]; i != c; i = D[i])

for(int i = U[c]; i != c; i = U[i])

for(int j = L[i]; j != i; j = L[j]) {
U[D[j]] = j;

for(int j = R[i]; j != i; j = R[j]) {
 U[D[j]] = U[j];

```
return;
int c = R[0];
for(int i = R[0]; i; i = R[i])
 if(S[i] < S[c]) c = i;
remove(c);
for(int i = D[c]; i != c; i = D[i]) {
  ans[d] = row[i]
  for(int j = R[i]; j != i; j = R[j])
    remove(col[j]);
  dance(d+1);
  for(int j = L[i]; j != i; j = L[j])
    resume(col[j]);
resume(c);
```

## 8.9 Tree Knapsack

```
int dp[N][K];PII obj[N];
vector<int> G[N];
void dfs(int u, int mx){
  for(int s: G[u]) {
    if(mx < obj[s].first) continue;</pre>
    for(int i=0;i<=mx-obj[s].FF;i++)</pre>
      dp[s][i] = dp[u][i];
    dfs(s, mx - obj[s].first);
    for(int i=obj[s].FF;i<=mx;i++)</pre>
      dp[u][i] = max(dp[u][i],
        dp[s][i - obj[s].FF] + obj[s].SS);
int main(){
  int n, k; cin >> n >> k;
  for(int i=1;i<=n;i++){</pre>
    int p; cin >> p;
    G[p].push_back(i);
    cin >> obj[i].FF >> obj[i].SS;
  dfs(0, k); int ans = 0;
  for(int i=0;i<=k;i++) ans = max(ans, dp[0][i]);</pre>
  cout << ans << '\n';
  return 0;
```

#### N Queens Problem

```
vector< int > solve( int n ) {
  // no solution when n=2. 3
  vector< int > ret;
  if ( n % 6 == 2 ) {
  for ( int i = 2 ; i <= n ; i += 2 )</pre>
       ret.push_back( i );
    ret.push_back( 3 ); ret.push_back( 1 ); for ( int i = 7 ; i <= n ; i += 2 )
       ret.push_back( i );
    ret.push_back( 5 );
  } else if ( n % 6 == 3 ) {
    for ( int i = 4 ; i <= n ; i += 2 )
       ret.push_back( i );
     ret.push_back( 2 );
    for ( int i = 5 ; i <= n ; i += 2 )
       ret.push_back( i );
     ret.push_back( 1 ); ret.push_back( 3 );
  } else {
    for ( int i = 2 ; i <= n ; i += 2 )
    ret.push_back( i );
for ( int i = 1 ; i <= n ; i += 2 )
       ret.push_back( i );
  return ret;
```