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5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.10 5.11 5.15 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.19 5.19 5.19 5.19 5.19 5.19 5.19	\begin{align*}{l} \text{ Enumeration} \text{ Strling Number} \text{.5.2.1 First Kind} \text{.5.2.2 Second Kind} \text{.3x+by=gcd} \text{.6x+by=gcd} .6x+by=	13 13 13 14 14 14 15 15 15 16 16 16 17 17 18 18
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	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22 6.0 6.1 6.2 6.3	\begin{align*}{l} \text{ Enumeration} \text{ Strling Number} \text{.} \text{.} \text{.} \text{ First Kind} \text{.}	13 13 13 14 14 14 15 15 15 16 16 16 16 17 17 18 18 18 19 19
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	5.1 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.22 5.23 6.4 6.2 6.3 6.4 6.5 6.6	\begin{align*}{l} \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \	133 13 13 14 14 14 14 15 15 15 16 16 16 16 16 17 18 18 18 18 19 19 19 19 19 19 19 19 19 19 19 19 19
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	5.1 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.22 5.23 6.4 6.2 6.3 6.4 6.5 6.6	\begin{align*}{l} \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \	133 13 13 14 14 14 14 15 15 15 16 16 16 16 16 17 18 18 18 18 19 19 19 19 19 19 19 19 19 19 19 19 19
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6.20 Circle Cover

7 St 7.1 7.2 7.3 7.4 7.5 7.6 7.5 7.5	2 Suffix Array 3 Suffix Automaton 4 KMP 5 Z value 5 Manacher 7 Lexico Smallest Rotation 8 BWT
8.4 8.3 8.3 8.4 8.5 8.6 8.8 8.8 8.8 8.8 8.8 8.8 8.8 8.8	Theorems 8.1.1 Kirchhoff's Theorem 8.1.2 Tutte's Matrix 8.1.3 Cayley's Formula 8.1.4 Erdős-Gallai theorem 8.1.5 Havel-Hakimi algorithm 8.1.6 Hall's marriage theorem 8.1.7 Euler's planar graph formula 8.1.8 Pick's theorem 8.1.9 Lucas's theorem 8.1.10 Matroid Intersection DP-opt Condition 8.2.1 totally monotone (concave/convex) 8.2.2 monge condition (concave/convex) 8.2.2 monge condition (concave/convex) 8.2.1 totally DP 4 Convex 1D/1D DP 5 Convex 1D/1D DP 6 Convex Hull Optimization 6 Josephus Problem 7 Tree Knapsack 8 N Queens Problem 8 Aliens Optimization 9 Aliens Optimization 10 Hilbert Curve
1 11	Basic vimro

```
se is nu ru et tgc sc hls cin cino+=j1 sw=4 sts=4 bs=2
    mouse=a encoding=utf-8 ls=2
syn on
colorscheme desert
filetype indent on
inoremap {<CR> {<CR>}<ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -std=c++17 -
     DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
     Wconversion -fsanitize=address,undefined -g && echo
      success<CR>
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 &&
      echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
```

1.2 Debug Macro

```
#ifdef KISEKI
 #define safe cerr<<__PRETTY_FUNCTION__\
<<" line "<<__LINE__<<" safe\n"</pre>
 #define debug(a...) qwerty(#a, a)
 #define orange(a...) dvorak(#a, a)
 using std::cerr;
 template <typename ...T>
void qwerty(const char *s, T ...a) {
  cerr << "\e[1;32m(" << s << ") = (";
  int cnt = sizeof...(T);</pre>
    (..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
 template <typename Iter>
 void dvorak(const char *s, Iter L, Iter R) {
  cerr << "\e[1;32m[ " << s << " ] = [ ";</pre>
   for (int f = 0; L != R; ++L)
  cerr << (f++ ? ", " : "") << *L;
cerr << " ]\e[0m\n";</pre>
 }
 #else
 #define safe ((void)0)
 #define debug(...) ((void)\theta)
#define orange(...) ((void)\theta)
 #endif
```

1.3 Increase Stack

```
const int size = 256 << 20;</pre>
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));
```

1.4 Pragma Optimization

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,sse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8000)
;
```

1.5 IO Optimization

```
| static inline int gc() {
    constexpr int B = 1<<20;
    static char buf[B], *p, *q;
    if(p == q &&
        (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
        return EOF;
    return *p++;
    }
    template < typename T >
    static inline bool gn( T &x ) {
        int c = gc(); T sgn = 1; x = 0;
        while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
        if(c == '-') sgn = -1, c = gc();
        if(c == EOF) return false;
        while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
        return x *= sgn, true;
}
```

2 Data Structure

2.1 Dark Magic

2.2 Link-Cut Tree

```
template <typename Val> class LCT {
private:
struct node
  int pa, ch[2];
 bool rev;
 Val v, v_prod, v_rprod;
 node() \; : \; pa\{0\}, \; ch\{0, \; 0\}, \; rev\{false\}, \; v\{\}, \; v\_prod\{\}, \\
    v_rprod{} {};
};
vector<node> nodes;
 set<pair<int, int>> edges;
bool is_root(int u) const {
 const int p = nodes[u].pa;
 return nodes[p].ch[0] != u and nodes[p].ch[1] != u;
bool is_rch(int u) const {
 return (not is_root(u)) and nodes[nodes[u].pa].ch[1]
    == u;
void down(int u) {
 if (auto &cnode = nodes[u]; cnode.rev) {
   if (cnode.ch[0]) set_rev(cnode.ch[0]);
   if (cnode.ch[1]) set_rev(cnode.ch[1]);
   cnode.rev = false;
 void up(int u) {
 auto &cnode = nodes[u];
 cnode.v_prod =
  nodes[cnode.ch[0]].v_prod * cnode.v * nodes[cnode.ch
    [1]].v_prod;
  cnode.v_rprod =
   nodes[cnode.ch[1]].v_rprod * cnode.v * nodes[cnode.
    ch[0]].v_rprod;
```

```
void set_rev(int u) {
  swap(nodes[u].ch[0], nodes[u].ch[1]);
  swap(nodes[u].v_prod, nodes[u].v_rprod);
 nodes[u].rev ^= 1;
 void rotate(int u) {
 int f = nodes[u].pa, g = nodes[f].pa, l = is_rch(u);
if (nodes[u].ch[l ^ 1])
  nodes[nodes[u].ch[1 ^ 1]].pa = f;
  if (not is_root(f))
   nodes[g].ch[is_rch(f)] = u;
  nodes[f].ch[1] = nodes[u].ch[1 ^ 1];
  nodes[u].ch[1 ^ 1] = f
  nodes[u].pa = g, nodes[f].pa = u;
 up(f);
 void splay(int u) {
  vector<int> stk = {u};
  while (not is_root(stk.back()))
  stk.push_back(nodes[stk.back()].pa);
  for (; not stk.empty(); stk.pop_back())
   down(stk.back())
  for(int f=nodes[u].pa;!is_root(u);f=nodes[u].pa){
   if(!is_root(f))rotate(is_rch(u)==is_rch(f)?f:u);
   rotate(u);
  }
 up(u);
 void access(int u) {
  int last = 0;
  for (int last = 0; u; last = u, u = nodes[u].pa) {
   splay(u);
   nodes[u].ch[1] = last;
   up(u);
 int find_root(int u) {
  access(u); splay(u);
  int la = 0:
  for (; u; la = u, u = nodes[u].ch[0]) down(u);
  return la;
 void change_root(int u) {
 access(u); splay(u); set_rev(u);
 void link(int x, int y) {
 change_root(y); nodes[y].pa = x;
 void split(int x, int y) {
 change_root(x); access(y); splay(y);
 void cut(int x, int y) {
 split(x, y)
  nodes[y].ch[0] = nodes[x].pa = 0;
 up(y);
public:
LCT(int n = 0) : nodes(n + 1) {}
 int add(const Val &v = {}) {
 nodes.push_back(v)
  return int(nodes.size()) - 2;
 int add(Val &&v) {
 nodes.emplace_back(move(v));
  return int(nodes.size()) - 2;
 void set_val(int u, const Val &v) {
 splay(++u); nodes[u].v = v; up(u);
 Val query(int x, int y) {
 split(++x, ++y)
  return nodes[y].v_prod;
 bool connected(int u, int v) { return find_root(++u)
    == find_root(++v); }
 void add_edge(int u, int v) {
 if (++u > ++v) swap(u, v)
  edges.emplace(u, v); link(u, v);
 void del_edge(int u, int v) {
 auto k = minmax(++u, ++v);
```

```
if (auto it = edges.find(k); it != edges.end()) {
                                                               } // sz(L) == k
   edges.erase(it); cut(u, v);
                                                               int getRank(node *o) { // 1-base
 }
                                                                int r = sz(o->lc) + 1;
};
                                                                for (;o->pa != nullptr; o = o->pa)
                                                                 if (o->pa->rc == o) r += sz(o->pa->lc) + 1;
2.3 LiChao Segment Tree
                                                                return r:
struct L {
int m, k, id;
L() : id(-1) {}
                                                               #undef sz
 L(int a, int b, int c) : m(a), k(b), id(c) {}
                                                              2.5 Linear Basis
 int at(int x) { return m * x + k; }
                                                             template <int BITS>
class LiChao {
                                                              struct LinearBasis {
                                                               array<uint64_t, BITS> basis;
private:
 int n; vector<L> nodes;
                                                               Basis() { basis.fill(0); }
 static int lc(int x) { return 2 * x + 1;
static int rc(int x) { return 2 * x + 2;
                                                               void add(uint64_t x)
                                                                for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
                                                                 if (basis[i] == 0) {
 void insert(int 1, int r, int id, L ln) {
  int m = (1 + r) >> 1;
                                                                  basis[i] = x;
  if (nodes[id].id == -1) {
                                                                  return:
   nodes[id] = ln;
                                                                 x ^= basis[i];
   return:
                                                               }
  bool atLeft = nodes[id].at(1) < ln.at(1);</pre>
  if (nodes[id].at(m) < ln.at(m)) {</pre>
                                                               bool ok(uint64_t x) {
  atLeft ^= 1
                                                               for (int i = 0; i < BITS; ++i)
   swap(nodes[id], ln);
                                                                 if ((x >> i) & 1) x ^= basis[i];
                                                                return x == 0;
  if (r - l == 1) return;
  if (atLeft) insert(l, m, lc(id), ln);
                                                             };
  else insert(m, r, rc(id), ln);
                                                                   Binary Search On Segment Tree
 int query(int 1, int r, int id, int x) {
                                                             // find_first = x -> minimal x s.t. check( [a, x) )
 int ret = 0, m = (1 + r) >> 1;
if (nodes[id].id != -1)
                                                             // find_last = x -> maximal x s.t. check( [x, b) )
                                                              template <typename C>
                                                             int find_first(int 1, const C &check) {
   ret = nodes[id].at(x);
  if (r - 1 == 1) return ret;
                                                              if (1 >= n) return n + 1;
  if (x < m) return max(ret, query(1, m, lc(id), x));
return max(ret, query(m, r, rc(id), x));</pre>
                                                               1 += sz;
                                                               for (int i = height; i > 0; i--)
                                                               propagate(1 >> i);
                                                               Monoid sum = identity;
public:
                                                               do {
 LiChao(int n_{-}) : n(n_{-}), nodes(n * 4) {}
                                                                while ((1 & 1) == 0) 1 >>= 1;
 void insert(L ln) { insert(0, n, 0, ln); }
                                                                if (check(f(sum, data[1]))) {
 int query(int x) { return query(0, n, 0, x); }
                                                                 while (1 < sz) {</pre>
                                                                  propagate(1);
                                                                  1 <<= 1;
2.4 Treap
                                                                  auto nxt = f(sum, data[1]);
namespace Treap{
                                                                  if (not check(nxt)) {
 #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                                   sum = nxt;
 struct node{
                                                                   1++:
                                                                  }
  int size:
  uint32_t pri;
 node *lc, *rc, *pa;
                                                                 return 1 + 1 - sz;
  node():size(0),pri(rand()),lc(0),rc(0),pa(0){}
  void pull() {
                                                                sum = f(sum, data[1++]);
   size = 1; pa = nullptr;
                                                               } while ((1 & -1) != 1);
   if ( lc ) { size += lc->size; lc->pa = this; }
                                                               return n + 1;
   if ( rc ) { size += rc->size; rc->pa = this; }
  }
                                                              template <typename C>
                                                              int find_last(int r, const C &check) {
                                                              if (r <= 0) return -1;</pre>
 node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
                                                               r += sz;
  if ( L->pri > R->pri ) {
                                                               for (int i = height; i > 0; i--)
  L->rc = merge( L->rc, R ); L->pull();
                                                               propagate((r - 1) >> i);
   return L;
                                                               Monoid sum = identity;
  } else {
                                                               do {
   R->lc = merge( L, R->lc ); R->pull();
   return R;
                                                                while (r > 1 and (r & 1)) r >>= 1;
                                                                if (check(f(data[r], sum))) {
  }
 }
                                                                 while (r < sz) {</pre>
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                                  propagate(r);
 if ( not rt ) L = R = nullptr;
                                                                  r = (r << 1) + 1;
                                                                  auto nxt = f(data[r], sum);
  else if( sz( rt->lc ) + 1 <= k ) {
                                                                  if (not check(nxt)) {
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                   sum = nxt:
   L->pull();
                                                                   r--:
  } else {
                                                                  }
   split_by_size( rt->lc, k, L, R->lc );
                                                                 return r - sz;
   R->pull();
```

```
sum = f(data[r], sum);
                                                                 dfs(v, u);
low[u] = min(low[u], low[v]);
 } while ((r & -r) != r);
return -1;
                                                                 if (low[v] > dfn[u]) bridge[t] = true;
                                                               }
3
    Graph
                                                              public:
                                                               void init(int n_) {
    2-SAT (SCC)
                                                                G.clear(); G.resize(n = n_);
class TwoSat{
                                                                low.assign(n, ecnt = 0);
private:
                                                                dfn.assign(n, 0);
int n;
vector<vector<int>> rG,G,sccs;
                                                               void add_edge(int u, int v) {
vector<int> ord,idx;
                                                                G[u].emplace_back(v, ecnt);
vector<bool> vis,result;
                                                                G[v].emplace_back(u, ecnt++);
void dfs(int u){
 vis[u]=true
                                                               void solve() {
  for(int v:G[u])
                                                                bridge.assign(ecnt, false);
  if(!vis[v]) dfs(v);
                                                                for (int i = 0; i < n; ++i)
 ord.push_back(u);
                                                                 if (not dfn[i]) dfs(i, i);
void rdfs(int u){
                                                               bool is_bridge(int x) { return bridge[x]; }
 vis[u]=false;idx[u]=sccs.size()-1;
                                                            } bcc_bridge;
  sccs.back().push_back(u);
                                                             3.3 BCC Vertex
  for(int v:rG[u])
  if(vis[v])rdfs(v);
                                                            class BCC_AP {
                                                              private:
public:
                                                               int n, ecnt;
 void init(int n_){
                                                               vector<vector<pair<int,int>>> G;
 G.clear();G.resize(n=n_);
                                                               vector<int> bcc, dfn, low, st;
  rG.clear();rG.resize(n);
                                                               vector<bool> ap, ins;
  sccs.clear();ord.clear();
                                                               void dfs(int u, int f) {
 idx.resize(n);result.resize(n);
                                                                dfn[u] = low[u] = dfn[f] + 1;
                                                                int ch = 0:
void add_edge(int u,int v){
                                                                for (auto [v, t]: G[u]) if (v != f) {
 G[u].push_back(v);rG[v].push_back(u);
                                                                 if (not ins[t]) {
                                                                  st.push_back(t);
void orr(int x,int y){
                                                                  ins[t] = true;
 if ((x^y)==1)return
 add_edge(x^1,y); add_edge(y^1,x);
                                                                 if (dfn[v]) {
                                                                  low[u] = min(low[u], dfn[v]);
bool solve(){
 vis.clear();vis.resize(n);
for(int i=0;i<n;++i)</pre>
                                                                 } ++ch; dfs(v, u)
                                                                 low[u] = min(low[u], low[v]);
   if(not vis[i])dfs(i);
                                                                 if (low[v] >= dfn[u]) {
  reverse(ord.begin(),ord.end());
                                                                  ap[u] = true;
  for (int u:ord){
                                                                  while (true) {
  if(!vis[u])continue;
                                                                   int eid = st.back(); st.pop_back();
   sccs.push_back(vector<int>());
                                                                   bcc[eid] = ecnt;
   rdfs(u);
                                                                   if (eid == t) break;
  for(int i=0;i<n;i+=2)</pre>
                                                                  ecnt++;
   if(idx[i]==idx[i+1])
    return false:
  vector<bool> c(sccs.size());
                                                                if (ch == 1 \text{ and } u == f) ap[u] = false;
  for(size_t i=0;i<sccs.size();++i){</pre>
   for(auto sij : sccs[i]){
                                                              public:
    result[sij]=c[i]
                                                               void init(int n_) {
    c[idx[sij^1]]=!c[i];
                                                                G.clear(); G.resize(n = n_);
                                                                ecnt = 0; ap.assign(n, false);
                                                                low.assign(n, 0); dfn.assign(n, 0);
 return true;
                                                               void add_edge(int u, int v) {
bool get(int x){return result[x];}
                                                               G[u].emplace_back(v, ecnt);
int get_id(int x){return idx[x];}
                                                                G[v].emplace_back(u, ecnt++);
int count(){return sccs.size();}
} sat2;
                                                               void solve() {
                                                                ins.assign(ecnt, false);
3.2 BCC Edge
                                                                bcc.resize(ecnt); ecnt = 0;
                                                                for (int i = 0; i < n; ++i)
if (not dfn[i]) dfs(i, i);</pre>
class BCC_Bridge {
private:
 int n, ecnt;
 vector<vector<pair<int,int>>> G;
                                                               int get_id(int x) { return bcc[x]; }
  vector<int> dfn, low;
                                                               int count() { return ecnt; }
  vector<br/>bool> bridge;
                                                               bool is_ap(int x) { return ap[x]; }
  void dfs(int u, int f) {
                                                            } bcc_ap;
   dfn[u] = low[u] = dfn[f] + 1;
                                                             3.4 Centroid Decomposition
   for (auto [v, t]: G[u]) {
    if (v == f) continue;
                                                            struct Centroid {
    if (dfn[v]) {
                                                              vector<vector<int64_t>> Dist;
     low[u] = min(low[u], dfn[v]);
                                                              vector<int> Parent, Depth;
                                                             vector<int64_t> Sub, Sub2;
     continue:
```

```
vector<int> Sz, Sz2;
Centroid(vector<vector<pair<int, int>>> g) {
 int N = g.size();
  vector<bool> Vis(N);
 vector < int > sz(N), mx(N);
 vector<int> Path;
 Dist.resize(N);
 Parent.resize(N);
 Depth.resize(N)
 auto DfsSz = [&](auto dfs, int x) -> void {
  Vis[x] = true; sz[x] = 1; mx[x] = 0;
   for (auto [u, w] : g[x]) {
    if (Vis[u]) continue;
    dfs(dfs, u)
    sz[x] += sz[u];
    mx[x] = max(mx[x], sz[u]);
   Path.push_back(x);
  auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
   -> void {
   Dist[x].push_back(D);Vis[x] = true;
   for (auto [u, w] : g[x]) {
   if (Vis[u]) continue;
    dfs(dfs, u, D + w);
  };
  auto Dfs = [&]
   (auto dfs, int x, int D = 0, int p = -1)->void {
Path.clear(); DfsSz(DfsSz, x);
   int M = Path.size();
   int C = -1;
   for (int u : Path) {
    if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
    Vis[u] = false;
   DfsDist(DfsDist, C);
   for (int u : Path) Vis[u] = false;
   Parent[C] = p; Vis[C] = true;
   Depth[C] = D;
   for (auto [u, w] : g[C]) {
   if (Vis[u]) continue;
dfs(dfs, u, D + 1, C);
  Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
  Sz.resize(N); Sz2.resize(N);
 void Mark(int v) {
 int x = v, z = -1;
  for (int i = Depth[v]; i >= 0; --i) {
   Sub[x] += Dist[v][i]; Sz[x]++;
   if (z != -1) {
    Sub2[z] += Dist[v][i];
    Sz2[z]++;
   z = x; x = Parent[x];
 }
int64_t Query(int v) {
 int64_t res = 0;
  int x = v, z = -1;
 for (int i = Depth[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
  if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
  z = x; x = Parent[x];
 }
  return res;
3.5 Directed Minimum Spanning Tree
```

```
struct DirectedMST { // find maximum
struct Edge {
 int u, v;
 int w:
 Edge(int u, int v, int w) : u(u), v(v), w(w) {}
vector<Edge> Edges;
void clear() { Edges.clear(); }
void addEdge(int a, int b, int w) { Edges.emplace_back
    (a, b, w); }
```

fa[x] = p;

return c ? p : val[x];

```
int solve(int root, int n) {
  vector<Edge> E = Edges;
  int ans = 0;
  while (true) {
   // find best in edge
   vector<int> in(n, -inf), prv(n, -1);
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
     in[e.v] = e.w;
     prv[e.v] = e.u;
   in[root] = 0;
   prv[root] = -1;
   for (int i = 0; i < n; i++)
    if (in[i] == -inf)
     return -inf;
   // find cycle
   int tot = 0;
   vector < int > id(n, -1), vis(n, -1);
   for (int i = 0; i < n; i++) {
    ans += in[i];
    for (int x = i; x != -1 && id[x] == -1; x = prv[x])
     if (vis[x] == i) {
      for (int y = prv[x]; y != x; y = prv[y])
       id[y] = tot;
      id[x] = tot++;
      break;
     vis[x] = i;
   if (!tot)
    return ans;
   for (int i = 0; i < n; i++)
    if (id[i] == -1)
     id[i] = tot++;
   // shrink
   for (auto &e : E) {
   if (id[e.u] != id[e.v])
     e.w -= in[e.v];
    e.u = id[e.u], e.v = id[e.v];
   n = tot:
   root = id[root];
  assert(false);
} DMST;
3.6 Dominator Tree
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
 // vertices are numbered from 0 to n - 1
 fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
 fill(fa, fa + n, -1); fill(val, val + n, -1);
 fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
 fill(dom, dom + n, -1); tk = 0;
for (int i = 0; i < n; ++i) {
  g[i].clear(); r[i].clear(); rdom[i].clear();
 }
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
 rev[dfn[x] = tk] = x;
 fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
 for (int u : g[x]) -
 if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
  r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
 if (fa[x] == x) return c ? -1 : x;
 int p = find(fa[x], 1);
 if (p == -1) return c ? fa[x] : val[x];
 if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
```

// chains : number of chain

```
// tl, tr[u] : subtree interval in the seq. of u
                                                                // head[i] : head of the chain i
vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
                                                                // chian[u] : chain id of the chain u is on
                                                                void predfs(int u, int f) {
  dep[u] = dep[pa[u] = f] + 1;
// p[i] = -2 if i is unreachable from s
dfs(s);
for (int i = tk - 1; i >= 0; --i)
                                                                 for (int v : G[u]) if (v != f) {
 for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
                                                                  predfs(v, u);
 if (i) rdom[sdom[i]].push_back(i);
                                                                  if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
  for (int &u : rdom[i]) {
                                                                   chain[u] = chain[v];
   int p = find(u);
   if (sdom[p] == i) dom[u] = i;
                                                                 if (chain[u] == 0) chain[u] = ++chains;
   else dom[u] = p;
                                                                void dfschain(int u, int f) {
  if (i) merge(i, rp[i]);
                                                                 tl[u] = timer++
}
                                                                 if (head[chain[u]] == -1)
vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)</pre>
                                                                  head[chain[u]] = u;
                                                                 for (int v : G[u])
  if (v != f and chain[v] == chain[u])
 if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
                                                                   dfschain(v, u);
                                                                 for (int v : G[u])
                                                                  if (v != f and chain[v] != chain[u])
                                                                   dfschain(v, u);
3.7 Edge Coloring
                                                                 tr[u] = timer;
// max(d_u) + 1 edge coloring, time: O(NM)
                                                               public:
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
                                                                LBD(int n) : timer(0), chains(0), G(n), tl(n), tr(n),
                                                                   chain(n), head(\hat{n}, -1), dep(\hat{n}), pa(\hat{n}) {}
for (int i = 0; i <= N; i++)
  for (int j = 0; j <= N; j++)
                                                                void add_edge(int u, int v) {
                                                                 G[u].push_back(v); G[v].push_back(u);
    C[i][j] = G[i][j] = 0;
                                                                void decompose() { predfs(0, 0); dfschain(0, 0); }
void solve(vector<pair<int, int>> &E, int N) {
int X[kN] = {}, a;
auto update = [&](int u)
                                                                PII get_subtree(int u) { return {tl[u], tr[u]}; }
                                                                vector<PII> get_path(int u, int v) {
                                                                 vector<PII> res;
 for (X[u] = 1; C[u][X[u]]; X[u]++);
                                                                 while (chain[u] != chain[v]) {
auto color = [&](int u, int v, int c) {
                                                                  if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
 int p = G[u][v];
                                                                   swap(u, v);
                                                                  int s = head[chain[u]];
 G[u][v] = G[v][u] = c;
  C[u][c] = v, C[v][c] = u;
                                                                  res.emplace_back(tl[s], tl[u] + 1);
 C[u][p] = C[v][p] = 0;
                                                                  u = pa[s];
  if (p) X[u] = X[v] = p;
 else update(u), update(v);
                                                                 if (dep[u] < dep[v]) swap(u, v);</pre>
                                                                 res.emplace_back(tl[v], tl[u] + 1);
 return p;
};
                                                                 return res;
auto flip = [&](int u, int c1, int c2) {
  int p = C[u][c1];
                                                               };
  swap(C[u][c1], C[u][c2]);
                                                                     Manhattan Minimum Spanning Tree
 if (p) G[u][p] = G[p][u] = c2;
  if (!C[u][c1]) X[u] = c1;
                                                               typedef Point<int> P;
 if (!C[u][c2]) X[u] = c2;
                                                               vector<array<int, 3>> manhattanMST(vector<P> ps) {
                                                                vi id(sz(ps));
  return p;
 };
                                                                iota(all(id), 0);
for (int i = 1; i <= N; i++) X[i] = 1;
                                                                vector<array<int, 3>> edges;
                                                                rep(k, 0, 4) {
  sort(all(id), [&](int i, int j) {
for (int t = 0; t < E.size(); t++) {</pre>
 auto [u, v] = E[t];
 int v\theta = v, c = X[u], c\theta = c, d;
                                                                  return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
  vector<pair<int, int>> L; int vst[kN] = {};
  while (!G[u][v0]) {
                                                                 map<int, int> sweep;
for (int i : id) {
   L.emplace_back(v, d = X[v]);
   if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
                                                                  for (auto it = sweep.lower_bound(-ps[i].y);
   c = color(u, L[a].first, c);
else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
                                                                     it != sweep.end(); sweep.erase(it++)) {
                                                                   int j = it->second;
     color(u, L[a].first, L[a].second);
                                                                   P d = ps[i] - ps[j];
   else if (vst[d]) break
                                                                   if (d.y > d.x) break;
   else vst[d] = 1, v = C[u][d];
                                                                   edges.push_back(\{d.y + d.x, i, j\});
  if (!G[u][v0]) {
                                                                  sweep[-ps[i].y] = i;
   for (; v; v = flip(v, c, d), swap(c, d));
   if (C[u][c0]) { a = int(L.size()) - 1;
                                                                 for (P &p : ps)
    while (--a >= 0 && L[a].second != c);
                                                                  if (k \& 1) p.x = -p.x;
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
                                                                  else swap(p.x, p.y);
   } else t--;
                                                                return edges; // [{w, i, j}, ...]
                                                                     MaxClique
3.8 Lowbit Decomposition
                                                               // contain a self loop \boldsymbol{u} to \boldsymbol{u}, than \boldsymbol{u} won't in clique
class LBD {
                                                               template < size_t MAXN >
int timer, chains;
                                                               class MaxClique{
vector<vector<int>> G;
                                                               private:
vector<int> tl, tr, chain, head, dep, pa;
                                                                using bits = bitset< MAXN >;
```

bits popped, G[MAXN], ans;

```
size_t deg[ MAXN ], deo[ MAXN ], n;
                                                                      if (k > mx) cs[++mx + 1].reset();
 void sort_by_degree() {
                                                                      cs[k][p] = 1;
  popped.reset();
                                                                      if (k < km) r[t++] = p;
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
    deg[ i ] = G[ i ].count();
                                                                     c.resize(m);
  for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;</pre>
                                                                     if (t) c[t - 1] = 0;
                                                                     for (int k = km; k <= mx; k++) {</pre>
                                                                      for (int p = int(cs[k]._Find_first());
    for ( size_t j = 0 ; j < n ; ++ j )
      if ( not popped[ j ] and deg[ j ] < mi )</pre>
                                                                           < kN; p = int(cs[k]._Find_next(p))) {
    mi = deg[ id = j ];
popped[ deo[ i ] = id ] = 1;
                                                                       r[t] = p; c[t++] = k;
    for( size_t u = G[ i ]._Find_first() ;
     u \( \) n ; u = G[ i ]._Find_next( u ) )
       -- deg[ u ];
                                                                    void dfs(vector<int> &r, vector<int> &c, int 1,
                                                                     bitset<kN> mask) {
  }
                                                                     while (!r.empty()) {
 void BK( bits R, bits P, bits X ) {
                                                                      int p = r.back(); r.pop_back();
 if (R.count()+P.count() <= ans.count()) return;</pre>
                                                                      mask[p] = 0;
  if ( not P.count() and not X.count() ) {
                                                                      if (q + c.back() <= ans) return;</pre>
   if ( R.count() > ans.count() ) ans = R;
                                                                      cur[q++] = p;
                                                                      vector<int> nr, nc;
   return:
                                                                      bitset<kN> nmask = mask & a[p];
  /* greedily chosse max degree as pivot
                                                                      for (int i : r)
  bits cur = P | X; size_t pivot = 0, sz = 0;
                                                                       if (a[p][i]) nr.push_back(i);
  for ( size_t u = cur._Find_first() ;
                                                                      if (!nr.empty()) {
   u < n ; u = cur._Find_next( u ) )
if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
                                                                       if (1 < 4) {
                                                                        for (int i : nr)
  cur = P & ( ~G[ pivot ] );
                                                                         d[i] = int((a[i] & nmask).count());
                                                                        sort(nr.begin(), nr.end(),
[&](int x, int y) {
  */ // or simply choose first
  bits cur = P & (~G[ ( P | X )._Find_first() ]);
  for ( size_t u = cur._Find_first()
                                                                           return d[x] > d[y];
   u < n ; u = cur._Find_next( u ) ) {
if ( R[ u ] ) continue;</pre>
                                                                         });
   R[u] = 1;
                                                                       csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
   BK(R, P \& G[u], X \& G[u]);
                                                                      } else if (q > ans) {
   R[u] = P[u] = 0, X[u] = 1;
                                                                       ans = q; copy(cur, cur + q, sol);
  }
                                                                      c.pop_back(); q--;
public:
                                                                     }
 void init( size_t n_ ) {
                                                                    int solve(bitset<kN> mask) { // vertex mask
  n = n_{-};
  for ( size_t i = 0 ; i < n ; ++ i )
G[ i ].reset();</pre>
                                                                     vector<int> r, c;
for (int i = 0; i < n; i++)</pre>
  ans.reset();
                                                                      if (mask[i]) r.push_back(i);
                                                                      for (int i = 0; i < n; i++)
d[i] = int((a[i] & mask).count());
                                                                     for
 void add_edges( int u, bits S ) { G[ u ] = S; }
 void add_edge( int u, int v ) {
                                                                     sort(r.begin(), r.end(),
 G[u][v] = G[v][u] = 1;
                                                                      [&](int i, int j) { return d[i] > d[j]; });
                                                                     csort(r, c);
 int solve() {
                                                                     dfs(r, c, 1, mask);
                                                                     return ans; // sol[0 ~ ans-1]
  sort_by_degree(); // or simply iota( deo... )
for ( size_t i = 0 ; i < n ; ++ i )</pre>
   deg[ i ] = G[ i ].count();
                                                                  } graph;
  bits pob, nob = 0; pob.set()
                                                                         Minimum Mean Cycle
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
  for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
                                                                  /* minimum mean cycle O(VE) */
   size_t v = deo[ i ];
                                                                   struct MMC{
   bits tmp; tmp[ v ] = 1;
                                                                   #define FZ(n) memset((n),0,sizeof(n))
   BK( tmp, pob & G[ v ], nob & G[ v ] );
                                                                   #define E 101010
   pob[ v ] = 0, nob[ v ] = 1;
                                                                   #define V 1021
                                                                   #define inf 1e9
  return static_cast< int >( ans.count() );
                                                                    struct Edge { int v,u; double c; };
                                                                    int n, m, prv[V][V], prve[V][V], vst[V];
                                                                    Edge e[E];
                                                                    vector<int> edgeID, cycle, rho;
3.11 MaxCliqueDyn
                                                                    double d[V][V];
                                                                    void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
constexpr int kN = 150;
struct MaxClique { // Maximum Clique
bitset<kN> a[kN], cs[kN];
                                                                    void add_edge( int vi , int ui , double ci )
 int ans, sol[kN], q, cur[kN], d[kN], n;
                                                                    { e[ m ++ ] = { vi , ui , ci }; }
void init(int _n) {
  n = _n, ans = q = 0;
                                                                    void bellman_ford() {
                                                                     for(int i=0; i<n; i++) d[0][i]=0;</pre>
  for (int i = 0; i < n; i++) a[i].reset();</pre>
                                                                     for(int i=0; i<n; i++) {</pre>
                                                                      fill(d[i+1], d[i+1]+n, inf);
for(int j=0; j<m; j++) {
  int v = e[j].v, u = e[j].u;
 void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
 void csort(vector<int> &r, vector<int> &c) {
                                                                       if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
  d[i+1][u] = d[i][v]+e[j].c;
 int mx = 1, km = max(ans - q + 1, 1), t = \hat{0},
    m = int(r.size())
  cs[1].reset(); cs[2].reset();
                                                                        prv[i+1][u] = v;
  for (int i = 0; i < m; i++) {
  int p = r[i], k = 1;</pre>
                                                                        prve[i+1][u] = j;
   while ((cs[k] & a[p]).count()) k++;
```

```
}
                                                                 int ans = INF;
                                                                 for (int i = 0; i < n; i++)</pre>
double solve(){
                                                                  ans = min(ans, dp[(1 << t) - 1][i]);
  // returns inf if no cycle, mmc otherwise
                                                                 return ans;
  double mmc=inf;
  int st = -1;
                                                              } solver;
  bellman_ford();
                                                                      Mo's Algorithm on Tree
  for(int i=0; i<n; i++) {</pre>
   double avg=-inf;
                                                              int q; vector< int > G[N];
   for(int k=0; k<n; k++) {</pre>
                                                              struct Que{
   if(d[n][i]<inf-eps]</pre>
                                                              int u, v, id;
} que[ N ];
     avg=max(avg,(d[n][i]-d[k][i])/(n-k));
    else avg=max(avg,inf);
                                                              int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
                                                               void dfs( int u, int f ) {
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
                                                                dfn[ u ] = dfn_++; int saved_rbp = stk_;
                                                                for ( int v : G[ u ] ) {
  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
                                                                 if ( v == f ) continue;
  for (int i=n; !vst[st]; st=prv[i--][st]) {
                                                                 dfs( v, u );
  vst[st]++;
                                                                 if ( stk_ - saved_rbp < SQRT_N ) continue;</pre>
   edgeID.PB(prve[i][st]);
                                                                 for ( ++ block_ ; stk_ != saved_rbp ; )
  block_id[ stk[ -- stk_ ] ] = block_;
  rho.PB(st);
 while (vst[st] != 2) {
                                                               stk[ stk_ ++ ] = u;
  int v = rho.back(); rho.pop_back();
   cycle.PB(v);
                                                               bool inPath[ N ];
   vst[v]++;
                                                               void Diff( int u ) {
 }
                                                               if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
 reverse(ALL(edgeID));
                                                                else { /*add this edge*/ }
 edgeID.resize(SZ(cycle));
  return mmc;
                                                               void traverse( int& origin_u, int u ) {
                                                               for ( int g = lca( origin_u, u )
} mmc;
                                                                origin_u != g ; origin_u = parent_of[ origin_u ] )
                                                                  Diff( origin_u );
3.13
      Minimum Steiner Tree
                                                                for (int v = u; v != origin_u; v = parent_of[v])
// Minimum Steiner Tree
                                                                Diff( v );
// 0(V 3^T + V^2 2^T)
                                                                origin_u = u;
struct SteinerTree {
#define V 33
                                                               void solve() {
#define T 8
                                                                dfs( 1, 1 );
#define INF 1023456789
                                                                while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
int n, dst[V][V], dp[1 << T][V], tdst[V];</pre>
                                                                sort( que, que + q, [](const Que& x, const Que& y) {
  return tie( block_id[ x.u ], dfn[ x.v ] )
void init(int _n) {
 n = _n;
                                                                     < tie( block_id[ y.u ], dfn[ y.v ] );
 for (int i = 0; i < n; i++) {
                                                                } );
  for (int j = 0; j < n; j++)
dst[i][j] = INF * (i != j);</pre>
                                                                int U = 1, V = 1;
                                                               for ( int i = 0 ; i < q ; ++ i ) {
  pass( U, que[ i ].u );
  pass( V, que[ i ].v );</pre>
void add_edge(int ui, int vi, int wi) {
                                                                 // we could get our answer of que[ i ].id
 dst[ui][vi] = min(dst[ui][vi], wi);
  dst[vi][ui] = min(dst[vi][ui], wi);
                                                              }
                                                               /*
void shortest_path() {
                                                              Method 2:
 for (int k = 0; k < n; k++)
                                                               dfs u:
  for (int i = 0; i < n; i++)</pre>
                                                               push u
   for (int j = 0; j < n; j++)
                                                                iterate subtree
     dst[i][j] = min(dst[i][j], dst[i][k] + dst[k][j]);
                                                               Let P = LCA(u, v), and St(u) <= St(v)
 int solve(const vector<int> &ter) {
                                                               if (P == u) query[St(u), St(v)]
 int t = (int)ter.size();
                                                               else query[Ed(u), St(v)], query[St(P), St(P)]
  for (int i = 1; i < (1 << t); i++)
  fill_n(dp[i], n, INF);
  fill_n(dp[0], n, 0);
                                                               3.15 Tree Hashing
  for (int msk = 1; msk < (1 << t); msk++) {</pre>
                                                              uint64_t hsah(int u, int f) {
   if (msk == (msk & (-msk))) {
                                                                uint64_t r = 127;
    int who = __lg(msk);
                                                                for (int v : G[ u ]) if (v != f) {
    for (int i = 0; i < n; i++)
                                                                 uint64_t hh = hsah(v, u);
     dp[msk][i] = dst[ter[who]][i];
                                                                 r=(r+(hh*hh)%1010101333)%1011820613;
    continue:
                                                                return r;
   for (int i = 0; i < n; i++)
    for (int submsk = (msk - 1) & msk; submsk; submsk =
     (submsk - 1) & msk)
                                                                     Virtural Tree
     dp[msk][i] = min(dp[msk][i], dp[submsk][i] + dp[
    msk ^ submsk][i]);
                                                               vector<pair<int, int>> build(vector<int> vs) { // tree
   for (int i = 0; i < n; i++) {
                                                                   0-base
    tdst[i] = INF;
                                                                vector<pair<int, int>> res;
    for (int j = 0; j < n; j++)
                                                                sort(vs.begin(), vs.end(), [](int i, int j) { return
     tdst[i] = min(tdst[i], dp[msk][j] + dst[j][i]);
                                                                   dfn[i] < dfn[j]; });
                                                                vector<int> s = {0};
   copy_n(tdst, n, dp[msk]);
                                                                for (int v : vs) if (v != 0) {
```

int o = lca(v, s.back());

```
if (o != s.back()) {
  while (s.size() >= 2 and dfn[s[s.size() - 2]] >= dfn
    [o]) {
    res.emplace_back(s[s.size() - 2], s.back());
    s.pop_back();
  }
  if (s.back() != o) {
    res.emplace_back(s.back(), o);
    s.back() = o;
  }
} s.push_back(v);
}
for (size_t i = 1; i < s.size(); ++i)
  res.emplace_back(s[i - 1], s[i]);
  return res;
}</pre>
```

4 Matching & Flow

4.1 Bipartite Matching

```
struct BipartiteMatching {
vector<int> X[N];
int fX[N], fY[N], n;
bitset<N> vis;
bool dfs(int x)
  for (auto i:X[x]) {
  if (vis[i]) continue;
   vis[i] = true;
   if (fY[i]==-1 || dfs(fY[i])){
    fY[fX[x] = i] = x;
    return true;
  return false;
void init(int n_, int m) {
 vis.reset();
 fill(X, X + (n = n_), vector<int>());
memset(fX, -1, sizeof(int) * n);
 memset(fY, -1, sizeof(int) * m);
 void add_edge(int x, int y){
 X[x].push_back(y); }
int solve() { // return how many pair matched
  int cnt = 0;
 for(int i=0;i<n;i++) {</pre>
  vis.reset();
   cnt += dfs(i);
  return cnt;
};
```

4.2 Dijkstra Cost Flow

```
// kN = #(vertices)
// MCMF.{Init, AddEdge, MincostMaxflow}
// MincostMaxflow(source, sink, flow_limit, &cost)
// => flow
using Pii = pair<int, int>
constexpr int kInf = 0x3f3f3f3f, kN = 500;
struct Edge {
int to, rev, cost, flow;
struct MCMF { // 0-based
int n{}, m{}, s{}, t{};
vector<Edge> graph[kN];
// Larger range for relabeling
int64_t dis[kN] = {}, h[kN] = {};
int p[kN] = {};
void Init(int nn) {
 n = nn;
 for (int i = 0; i < n; i++) graph[i].clear();</pre>
void AddEdge(int u, int v, int f, int c) {
 graph[u].push_back({v
   static_cast<int>(graph[v].size()), c, f});
  graph[v].push_back(
   {u, static_cast<int>(graph[u].size()) - 1,
    -c, 0});
```

```
bool Dijkstra(int &max_flow, int64_t &cost) {
  priority_queue<Pii, vector<Pii>, greater
pq;
  fill_n(dis, n, kInf);
  dis[s] = 0;
  pq.emplace(0, s);
  while (!pq.empty()) {
   auto u = pq.top();
   pq.pop();
   int v = u.second;
   if (dis[v] < u.first) continue;</pre>
   for (auto &e : graph[v]) {
    auto new_dis =
    dis[v] + e.cost + h[v] - h[e.to];
if (e.flow > 0 && dis[e.to] > new_dis) {
     dis[e.to] = new_dis;
     p[e.to] = e.rev;
     pq.emplace(dis[e.to], e.to);
   }
  if (dis[t] == kInf) return false;
  for (int i = 0; i < n; i++) h[i] += dis[i];</pre>
  int d = max_flow;
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
d = min(d, graph[e.to][e.rev].flow);
  max_flow -= d;
  cost += int64_t(d) * h[t];
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   e.flow += d;
   graph[e.to][e.rev].flow -= d;
  return true;
 int MincostMaxflow(
  int ss, int tt, int max_flow, int64_t &cost) {
  this->s = ss, this->t = tt;
  cost = 0;
  fill_n(h, n, 0);
  auto orig_max_flow = max_flow;
  while (Dijkstra(max_flow, cost) && max_flow) {}
  return orig_max_flow - max_flow;
};
4.3 Dinic
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
    int to, rev;
    Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
      int u = bfs.front(); bfs.pop();
       for (auto e: G[u]) {
         if (e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
         bfs.push(e.to); lv[e.to] = lv[u] + 1;
      }
    return lv[ed] != -1;
  Cap DFS(int u, Cap f){
  if (u == ed) return f;
    Cap ret = 0;
    for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
      auto &e = G[u][i];
      if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
      Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
```

G[e.to][e.rev].cap += nf;

```
if (f == 0) return ret;
     if (ret == 0) lv[u] = -1;
     return ret;
public:
   void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
  G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
  st = st_, ed = ed_; Cap ret = 0;
     while (BFS()) {
        idx.assign(n, 0);
        Cap f = DFS(st, numeric_limits<Cap>::max());
        ret += f;
        if (f == 0) break;
     return ret;
};
        Flow Models
```

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}$.

 - For each edge (x, y, l, u), connect x → y with capacity u − l.
 For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t\to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T.If $f
 eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the
 - maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect t o s with capacity ∞ and let the flow from S to T be f'. If $f+f'\neq \sum_{v\in V, in(v)>0}in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipartite graph(X, Y)
 - $\begin{array}{l} \text{ If } (X,Y) \\ \text{ 1. Redirect every edge: } y \to x \text{ if } (x,y) \in M, x \to y \text{ otherwise.} \\ \text{ 2. DFS from unmatched vertices in } X. \\ \text{ 3. } x \in X \text{ is chosen iff } x \text{ is unvisited.} \\ \text{ 4. } y \in Y \text{ is chosen iff } y \text{ is visited.} \end{array}$
- · Minimum cost cyclic flow

 - 1. Consruct super source S and sink T 2. For each edge (x,y,c), connect $x\to y$ with (cost,cap)=(c,1) if
 - c>0, otherwise connect $y\to x$ with (cost,cap)=(-c,1)3. For each edge with c<0, sum these cost as K, then increase d(y)

 - by 1, decrease d(x) by 1 4. For each vertex v with d(v) > 0, connect $S \to v$ with (cost, cap) =
 - 5. For each vertex v with d(v)<0 , connect $v\to T$ with (cost,cap)=(0,-d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights 3. Connect source $s \to v, v \in G$ with capacity K

 - 4. For each edge (u,v,w) in G, connect u o v and v o u with capacity
 - 5. For $v \in \mathit{G}$, connect it with sink $v \to t$ with capacity K + 2T - $(\sum_{e \in E(v)} w(e)) - 2w(v)$ 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge
 - (v,t) with capacity $-p_v$. 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_x x + \sum_{y} c_y \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with ca-
- pacity c_y .

 2. Create edge (x,y) with capacity c_{xy} .

 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

```
4.5 General Graph Matching
```

```
namespace matching {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
  for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
 for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
 g[u].push_back(v);
 g[v].push_back(u);
int Find(int u) {
 return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
 static int tk = 0; tk++;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
void Blossom(int x, int y, int 1) {
 while (Find(x) != 1) {
 pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
if (fa[x] == x) fa[x] = 1;
  if (fa[y] == y) fa[y] = 1;
  x = pre[y];
bool Bfs(int r, int n) {
 for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;
 while (!q.empty()) q.pop();
 q.push(r);
 s[r] = 0;
 while (!q.empty()) {
  int x = q.front(); q.pop();
  for (int u : g[x]) {
  if (s[u] == -1) {
    pre[u] = x, s[u] = 1;
    if (match[u] == n) {
     for (int a = u, b = x, last; b != n; a = last, b =
     pre[a])
      last = match[b], match[b] = a, match[a] = b;
     return true;
    q.push(match[u]);
    s[match[u]] = 0;
   } else if (!s[u] && Find(u) != Find(x)) {
    int 1 = LCA(u, x, n);
    Blossom(x, u, 1);
    Blossom(u, x, 1)
  }
 return false;
int Solve(int n) {
 int res = 0;
 for (int x = 0; x < n; ++x) {
  if (match[x] == n) res += Bfs(x, n);
 }
 return res;
    Global Min-Cut
```

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
pair<int, int> phase(int n) {
 memset(v, false, sizeof(v));
 memset(g, 0, sizeof(g));
```

if (!v1[x] && d > s1k[x]) d = s1k[x];

```
int s = -1, t = -1;
                                                                   for (int x = 0; x < n; ++x) {
 while (true) {
                                                                    if (v1[x]) h1[x] += d;
                                                                    else slk[x] -= d;
  int c = -1;
  for (int i = 0; i < n; ++i) {</pre>
                                                                    if (vr[x]) hr[x] -= d;
   if (del[i] || v[i]) continue;
                                                                   for (int x = 0; x < n; ++x)
if (!v1[x] && !slk[x] && !check(x)) return;</pre>
   if (c == -1 || g[i] > g[c]) c = i;
  if (c == -1) break;
 v[s = t, t = c] = true;
for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
                                                               public:
                                                                void init( int n_ ) {
   g[i] += w[c][i];
                                                                 qu.resize(n = n_{-});
                                                                 fl.assign(n, -1); fr.assign(n, -1);
hr.assign(n, 0); hl.resize(n);
  }
                                                                 w.assign(n, vector<lld>(n));
 return make_pair(s, t);
                                                                 slk.resize(n); pre.resize(n);
int mincut(int n) {
                                                                 vl.resize(n); vr.resize(n);
 int cut = 1e9;
 memset(del, false, sizeof(del));
                                                                void set_edge( int u, int v, 1ld x ) {w[u][v] = x;}
 for (int i = 0; i < n - 1; ++i) {
                                                                1ld solve() {
  int s, t; tie(s, t) = phase(n);
                                                                 for (int i = 0; i < n; ++i)
                                                                  hl[i] = *max_element(w[i].begin(), w[i].end());
  del[t] = true; cut = min(cut, g[t]);
  for (int j = 0; j < n; ++j) {
                                                                 for (int i = 0; i < n; ++i) bfs(i);</pre>
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
                                                                 11d res = 0:
  }
                                                                 for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
                                                                 return res:
                                                                }
 return cut;
                                                               } km;
                                                               4.9 Minimum Cost Circulation
4.7
     GomoryHu Tree
                                                               struct Edge { int to, cap, rev, cost; };
int g[maxn];
                                                               vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
vector<edge> GomoryHu(int n){
 vector<edge> rt;
 for(int i=1;i<=n;++i)g[i]=1;</pre>
                                                               bool mark[kN];
 for(int i=2;i<=n;++i){</pre>
                                                               int NegativeCycle(int n) {
                                                                memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
  int t=g[i];
  flow.reset(); // clear flows on all edge
  rt.push_back({i,t,flow(i,t)});
                                                                int upd = -1;
 flow.walk(i); // bfs points that connected to i (use
  edges not fully flow)
                                                                for (int i = 0; i <= n; ++i)
                                                                 for (int j = 0; j < n; ++j) {
  for(int j=i+1;j<=n;++j){</pre>
                                                                   int idx = 0;
                                                                   for (auto &e : g[j]) {
   if(g[j]==t && flow.connect(j))g[j]=i; // check if i
                                                                    if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
    can reach j
 }
                                                                     dist[e.to] = dist[j] + e.cost;
                                                                     pv[e.to] = j, ed[e.to] = idx;
 return rt;
                                                                     if (i == n) {
                                                                      upd = j;
                                                                      while(!mark[upd])mark[upd]=1,upd=pv[upd];
4.8 Kuhn Munkres
                                                                      return upd;
class KM {
                                                                     }
private:
 static constexpr 1ld INF = 1LL << 60;</pre>
                                                                    idx++;
 vector<lld> hl,hr,slk;
 vector<int> fl,fr,pre,qu;
                                                                 }
 vector<vector<lld>> w;
 vector<bool> v1,vr;
                                                                return -1;
 int n, ql, qr;
 bool check(int x) {
                                                               int Solve(int n) {
  if (v1[x] = true, f1[x] != -1)
                                                                int rt = -1, ans = 0;
   return vr[qu[qr++] = f1[x]] = true;
                                                                while ((rt = NegativeCycle(n)) >= 0) {
  while (x != -1) swap(x, fr[fl[x] = pre[x]]);
                                                                 memset(mark, false, sizeof(mark));
                                                                 vector<pair<int, int>> cyc;
  return false;
                                                                 while (!mark[rt]) {
 void bfs(int s) {
                                                                   cyc.emplace_back(pv[rt], ed[rt]);
  fill(slk.begin(), slk.end(), INF);
                                                                   mark[rt] = true;
  fill(v1.begin(), v1.end(), false);
                                                                   rt = pv[rt];
  fill(vr.begin(), vr.end(), false);
  ql = qr = 0;
                                                                 reverse(cyc.begin(), cyc.end());
  vr[qu[qr++] = s] = true;
                                                                 int cap = kInf;
  while (true) {
                                                                 for (auto &i : cyc)
                                                                   auto &e = g[i.first][i.second];
   11d d;
   while (ql < qr) {</pre>
                                                                   cap = min(cap, e.cap);
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
     if(!v1[x]\&s1k[x]>=(d=h1[x]+hr[y]-w[x][y])){
                                                                 for (auto &i : cyc) {
                                                                   auto &e = g[i.first][i.second];
      if (pre[x] = y, d) slk[x] = d;
      else if (!check(x)) return;
                                                                   e.cap -= cap;
     }
                                                                   g[e.to][e.rev].cap += cap;
    }
                                                                   ans += e.cost * cap;
   d = INF;
   for (int x = 0; x < n; ++x)
                                                                return ans;
```

4.10 Minimum Cost Maximum Flow

```
class MiniCostMaxiFlow{
 using Cap = int; using Wei = int64_t;
 using PCW = pair<Cap,Wei>;
 static constexpr Cap INF_CAP = 1 << 30;</pre>
 static constexpr Wei INF_WEI = 1LL<<60;</pre>
private:
 struct Edge{
  int to, back;
Cap cap; Wei wei;
  Edge() {}
 Edge(int a,int b, Cap c, Wei d):
   to(a),back(b),cap(c),wei(d) {}
 };
 int ori, edd;
 vector<vector<Edge>> G;
 vector<int> fa, wh;
 vector<bool> inq;
 vector<Wei> dis;
 PCW SPFA(){
  fill(inq.begin(),inq.end(),false)
  fill(dis.begin(), dis.end(), INF_WEI);
  queue<int> qq; qq.push(ori);
  dis[ori] = 0;
  while(not qq.empty()){
   int u=qq.front();qq.pop();
   inq[u] = false;
   for(int i=0;i<SZ(G[u]);++i){</pre>
    Edge e=G[u][i];
    int v=e.to; Wei d=e.wei;
    if(e.cap <= 0 | |dis[v] <= dis[u] + d)
     continue
    dis[v] = dis[u] + d;
    fa[v] = u, wh[v] = i;
if (inq[v]) continue;
    qq.push(v);
    inq[v] = true;
  if(dis[edd]==INF_WEI) return {-1, -1};
  Cap mw=INF_CAP;
  for(int i=edd;i!=ori;i=fa[i])
   mw=min(mw,G[fa[i]][wh[i]].cap);
  for (int i=edd;i!=ori;i=fa[i]){
   auto &eg=G[fa[i]][wh[i]];
   eg.cap -= mw;
   G[eg.to][eg.back].cap+=mw;
  return {mw, dis[edd]};
public:
 void init(int n){
  G.clear();G.resize(n);
  fa.resize(n);wh.resize(n);
  inq.resize(n); dis.resize(n);
 void add_edge(int st, int ed, Cap c, Wei w){
  G[st].emplace_back(ed,SZ(G[ed]),c,w);
  G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
 PCW solve(int a, int b){
 ori = a, edd = b;
  Cap cc=0; Wei ww=0;
  while(true)
  PCW ret=SPFA();
   if(ret.first==-1) break;
   cc+=ret.first;
   ww+=ret.first * ret.second;
  }
  return {cc,ww};
 }
} mcmf;
```

4.11 Maximum Weight Graph Matching

```
struct WeightGraph {
  static const int inf = INT_MAX;
  static const int maxn = 514;
  struct edge {
  int u, v, w;
  edge(){}
  edge(int u, int v, int w): u(u), v(v), w(w) {}
```

```
int n, n_x;
edge g[maxn * 2][maxn * 2];
int lab[maxn * 2];
int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa
   [maxn * 2];
int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
   maxn * 21
vector<int> flo[maxn * 2];
queue<int> q;
int e_delta(const edge &e) { return lab[e.u] + lab[e.v
   ] - g[e.u][e.v].w * 2; }
void update_slack(int u, int x) { if (!slack[x] ||
   e_delta(g[u][x]) < e_delta(g[slack[x]][x])) slack[x]
   ] = u; }
void set_slack(int x) {
 slack[x] = 0;
 for (int u = 1; u <= n; ++u)
  if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
   update_slack(u, x);
void q_push(int x) {
if (x \le n) q.push(x);
 else for (size_t i = 0; i < flo[x].size(); i++)</pre>
   q_push(flo[x][i]);
void set_st(int x, int b) {
st[x] = b;
 if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)</pre>
    set_st(flo[x][i], b);
int get_pr(int b, int xr) {
 int pr = find(flo[b].begin(), flo[b].end(), xr) - flo
   [b].begin();
 if (pr % 2 == 1) {
  reverse(flo[b].begin() + 1, flo[b].end());
 return (int)flo[b].size() - pr;
}
 return pr;
void set_match(int u, int v) {
match[u] = g[u][v].v;
 if (u <= n) return;</pre>
 edge e = g[u][v];
 int xr = flo_from[u][e.u], pr = get_pr(u, xr)
 for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo</pre>
   [u][i ^ 1]);
set_match(xr, v);
rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
   end());
void augment(int u, int v) {
for (; ; ) {
  int xnv = st[match[u]];
  set_match(u, v);
 if (!xnv) return
  set_match(xnv, st[pa[xnv]]);
  u = st[pa[xnv]], v = xnv;
int get_lca(int u, int v) {
 static int t = 0;
 for (++t; u || v; swap(u, v)) {
 if (u == 0) continue;
if (vis[u] == t) return u;
 vis[u] = t;
 u = st[match[u]];
 if (u) u = st[pa[u]];
 return 0;
void add_blossom(int u, int lca, int v) {
 int b = n + 1;
 while (b <= n_x && st[b]) ++b;</pre>
 if (b > n_x) ++n_x;
 lab[b] = 0, S[b] = 0;
 match[b] = match[lca];
 flo[b].clear();
 flo[b].push_back(lca);
 for (int x = u, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
```

]]), q_push(y);

```
reverse(flo[b].begin() + 1, flo[b].end())
 for (int x = v, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
 set_st(b, b);
 for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].w
   = 0;
 for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
 for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  for (int x = 1; x <= n_x; ++x)
   if (g[b][x].w == 0 \mid \mid e_delta(g[xs][x]) < e_delta(g[xs][x])
   [b][x]))
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
  for (int x = 1; x <= n; ++x)
   if (flo_from[xs][x]) flo_from[b][x] = xs;
 set_slack(b);
}
void expand_blossom(int b) {
 for (size_t i = 0; i < flo[b].size(); ++i)
set_st(flo[b][i], flo[b][i]);</pre>
 int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
   xr);
 for (int i = 0; i < pr; i += 2) {
  int xs = flo[b][i], xns = flo[b][i + 1];</pre>
  pa[xs] = g[xns][xs].u;
  S[xs] = 1, S[xns] = 0;
  slack[xs] = 0, set_slack(xns);
  q_push(xns);
 S[xr] = 1, pa[xr] = pa[b];
 for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  S[xs] = -1, set_slack(xs);
 st[b] = 0;
bool on_found_edge(const edge &e) {
 int u = st[e.u], v = st[e.v];
 if (S[v] == -1) {
  pa[v] = e.u, S[v] = 1;
  int nu = st[match[v]];
  slack[v] = slack[nu] = 0;
 S[nu] = 0, q_push(nu);
} else if (S[v] == 0) {
  int lca = get_lca(u, v);
  if (!lca) return augment(u,v), augment(v,u), true;
  else add_blossom(u, lca, v);
 return false;
bool matching() {
 memset(S + 1, -1, sizeof(int) * n_x;
memset(slack + 1, 0, sizeof(int) * n_x);
 q = queue<int>();
 for (int x = 1; x <= n_x; ++x)
 if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0,
   q_push(x);
 if (q.empty()) return false;
 for (; ; ) {
  while (q.size()) {
   int u = q.front(); q.pop();
   if (S[st[u]] == 1) continue;
   for (int v = 1; v <= n; ++v)
    if (g[u][v].w > 0 && st[u] != st[v]) {
     if (e_delta(g[u][v]) == 0) {
      if (on_found_edge(g[u][v])) return true;
     } else update_slack(u, st[v]);
    }
  int d = inf;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b && S[b] == 1) d = min(d, lab[b] / 2)
  for (int x = 1; x <= n_x; ++x)
   if (st[x] == x && slack[x]) {
  if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x])
    else if (S[x] == 0) d = min(d, e_delta(g[slack[x
   ]][x]) / 2);
```

```
for (int u = 1; u <= n; ++u) {
    if (S[st[u]] == 0) {
     if (lab[u] <= d) return 0;</pre>
      lab[u] -= d;
    } else if (S[st[u]] == 1) lab[u] += d;
   for (int b = n + 1; b \le n_x; ++b)
    if (st[b] == b) {
      if (S[st[b]] == 0) lab[b] += d * 2;
     else if (S[st[b]] == 1) lab[b] -= d * 2;
    q = queue<int>();
   for (int x = 1; x <= n_x; ++x)
if (st[x] == x && slack[x] && st[slack[x]] != x &&
     e_delta(g[slack[x]][x]) == 0)
     if (on_found_edge(g[slack[x]][x])) return true;
    for (int b = n + 1; b <= n_x; ++b)
    if (st[b] == b && S[b] == 1 && lab[b] == 0)
     expand_blossom(b);
  return false;
 pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n_x = n:
  int n_matches = 0;
  long long tot_weight = 0;
  for (int u = 0; u \le n; t+u) st[u] = u, flo[u].clear
     ();
  int w_max = 0;
  for (int u = 1; u <= n; ++u)
   for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);</pre>
    w_{max} = max(w_{max}, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
  while (matching()) ++n_matches;
  for (int u = 1; u <= n; ++u)
   if (match[u] && match[u] < u)</pre>
    tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
 void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
     g[vi][ui].w = wi; }
 void init(int _n) {
  n = _n;
  for (int u = 1; u <= n; ++u)
   for (int v = 1; v <= n; ++v)
     g[u][v] = edge(u, v, 0);
};
```

5 Math

5.1 $\lfloor \frac{n}{i} \rfloor$ Enumeration

 $T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor$

5.2 Strling Number

5.2.1 First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1) \dots (x+n-1) = \sum_{k=0}^n S_1(n,k) x^k$$

$$g(x) = x(x+1) \dots (x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.2.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

```
S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}
```

5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 \pmod{n}
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
if (y == 0) g=x, a=1, b=0;
else exgcd(y, x\%y, g, b, a), b=(x/y)*a;
```

5.4 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
vector<T> d(output.size() + 1), me, he;
for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
  for (size_t j = 0; j < me.size(); ++j)</pre>
  d[i] += output[i - j - 2] * me[j];
if ((d[i] -= output[i - 1]) == 0) continue;
  if (me.empty()) {
   me.resize(f = i);
   continue;
 vector<T> o(i - f - 1);
T k = -d[i] / d[f]; o.push_back(-k);
 for (T x : he) o.push_back(x * k);
 if (o.size() < me.size()) o.resize(me.size());</pre>
  for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
  if (i-f+he.size() >= me.size()) he = me, f = i;
 me = o:
return me;
```

5.5 Charateristic Polynomial

```
vector<vector<int>> Hessenberg(const vector<vector<int</pre>
    >> &A) {
 int N = A.size();
vector<vector<int>> H = A;
for (int i = 0; i < N - 2; ++i) {
 if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {
    if (H[j][i]) {
     for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
    ][k]);
     for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k
    ][j]);
    }
   }
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
 for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
   * H[i + 1][k] * (kP - coef)) % kP;
for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
return H:
vector<int> CharacteristicPoly(const vector<vector<int
    >> &A) {
int N = A.size();
 auto H = Hessenberg(A);
for (int i = 0; i < N; ++i) {
 for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
vector<vector<int>> P(N + 1, vector<int>(N + 1));
P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
 P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
    1];
  int val = 1;
 for (int j = i - 1; j >= 0; --j) {
  int coef = 1LL * val * H[j][i - 1] % kP;
   for (int k = 0; k \le j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
```

```
if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 if (N & 1) {
  for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
 return P[N];
5.6 Chinese Remainder
1ld crt(lld ans[], lld pri[], int n){
 lld M = 1, ret = 0;
 for(int i=0;i<n;i++) M *= pri[i];</pre>
 for(int i=0;i<n;i++){</pre>
  lld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
  ret += (ans[i]*(M/pri[i])%M * iv)%M;
  ret %= M;
 return ret;
}
/*
Another:
x = a1 \% m1
x = a2 \% m2
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
0 <= x < lcm(m1, m2)
5.7 De-Bruijn
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
 if (t > n) {
  if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
 } else {
  aux[t] = aux[t - p];
  db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {
   aux[t] = i;
   db(t + 1, t, n, k);
  }
int de_bruijn(int k, int n) {
 // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
 if (k == 1) {
  res[0] = 0;
  return 1:
 for (int i = 0; i < k * n; i++) aux[i] = 0;
 sz = 0;
 db(1, 1, n, k);
 return sz;
5.8 DiscreteLog
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
  // x^? \equiv y (mod M)
  Int t = 1, c = 0, g = 1;
  for (Int M_ = M; M_ > 0; M_ >>= 1)
    g = g * x % M;
  for (g = gcd(g, M); t % g != 0; ++c) {
    if (t == y) return c;
    t = t * x % M;
  if (y % g != 0) return -1;
  t /= g, y /= g, M /= g;
  Int h = 0, gs = 1;
for (; h * h < M; ++h) gs = gs * x % M;
  unordered_map<Int, Int> bs;
  for (Int s = 0; s < h; bs[y] = ++s)
    y = y * x % M;
  for (Int s = 0; s < M; s += h) {
```

t = t * gs % M;

if (bs.count(t)) return c + s + h - bs[t];

5.10 ExtendedFloorSum

```
\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \\ - \frac{\frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1)}{-h(c, c-b-1, a, m-1)),} & \text{otherwise} \end{cases} \\ h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{c} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}
```

5.11 Fast Fourier Transform

```
const int mod = 1000000007:
const int M1 = 985661441; // G = 3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
  static_assert (M1 <= M2 && M2 <= M3);
 constexpr int64_t r12 = modpow(M1, M2-2, M2);
 constexpr int64_t r13 = modpow(M1, M3-2, M3);
 constexpr int64_t r23 = modpow(M2, M3-2, M3);
 constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
 B = (B - A + M2) * r12 % M2;
 C = (C - A + M3) * r13 % M3;
 C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
for (int i = 0; i <= maxn; i++)</pre>
 omega[i] = cplx(cos(2 * pi * j / maxn),
     sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
int z = __builtin_ctz(n) - 1;
for (int i = 0; i < n; ++i) {
 int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^{=(i >> j & 1)<<(z - j);
 if (x > i) swap(v[x], v[i]);
for (int s = 2; s <= n; s <<= 1) {
 int z = s >> 1;
for (int i = 0; i < n; i += s) {</pre>
  for (int k = 0; k < z; ++k) {
  cplx x = v[i + z + k] * omega[maxn / s * k];
    v[i + z + k] = v[i + k] - x;
    v[i + k] = v[i + k] + x;
void ifft(vector<cplx> &v, int n) {
fft(v, n); reverse(v.begin() + 1, v.end());
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);
VL convolution(const VI &a, const VI &b) {
// Should be able to handle N <= 10^5, C <= 10^4
```

```
int sz = 1;
 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 vector<cplx> v(sz);
 for (int i = 0; i < sz; ++i) {
  double re = i < a.size() ? a[i] : 0;</pre>
  double im = i < b.size() ? b[i] : 0;</pre>
  v[i] = cplx(re, im);
 fft(v, sz);
 for (int i = 0; i <= sz / 2; ++i) {
int j = (sz - i) & (sz - 1);
  cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
  * cplx(0, -0.25);
if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
     ].conj()) * cplx(0, -0.25);
  v[i] = x;
 ifft(v, sz);
 VL c(sz);
 for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
 return c:
VI convolution_mod(const VI &a, const VI &b, int p) {
 int sz = 1:
 while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
 vector<cplx> fa(sz), fb(sz);
for (int i = 0; i < (int)a.size(); ++i)</pre>
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 for (int i = 0; i < (int)b.size(); ++i)
fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
 fft(fa, sz), fft(fb, sz);
 double r = 0.25 / sz;
 cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
 for (int i = 0; i <= (sz >> 1); ++i) {
  int j = (sz - i) & (sz - 1);
  cplx a1 = (fa[i] + fa[j].conj());
  cplx a2 = (fa[i] - fa[j].conj()) * r2;
  cplx b1 = (fb[i] + fb[j].conj()) * r3;
  cplx b2 = (fb[i] - fb[j].conj()) * r4;
  if (i != j) {
   cplx c1 = (fa[j] + fa[i].conj());

cplx c2 = (fa[j] - fa[i].conj()) * r2;

cplx d1 = (fb[j] + fb[i].conj()) * r3;
   cplx d2 = (fb[j] - fb[i].conj()) * r4;
   fa[i] = c1 * d1 + c2 * d2 * r5;
fb[i] = c1 * d2 + c2 * d1;
  fa[j] = a1 * b1 + a2 * b2 * r5;
  fb[j] = a1 * b2 + a2 * b1;
 fft(fa, sz), fft(fb, sz);
 vector<int> res(sz);
 for (int i = 0; i < sz; ++i) {
  long long a = round(fa[i].re), b = round(fb[i].re),
        c = round(fa[i].im);
  res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
 }
 return res;
}}
5.12 FloorSum
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
// @return sum_\{i=0\}^{n-1} floor((ai + b)/m) mod 2^64
llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
 llu ans = 0;
 while (true) {
  if (a >= m) {
   ans += n * (n - 1) / 2 * (a / m); a %= m;
  if (b >= m) {
   ans += n * (b / m); b %= m;
  llu y_max = a * n + b;
  if (y_max < m) break;</pre>
  // y_{max} < m * (n + 1)
  // floor(y_max / m) <= n
  n = (11u)(y_max / m), b = (11u)(y_max % m);
  swap(m, a);
 return ans;
```

11u odd = (x - 1) >> t;

```
11d floor_sum(11d n, 11d m, 11d a, 11d b) {
                                                                for (llu m:
                                                                 {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
llu ans = 0;
if (a < 0) {
                                                                 if (witn(mpow(m % x, odd, x), odd, x, t))
 11u \ a2 = (a \% m + m) \% m;
                                                                  return false:
 ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
                                                                return true:
                                                               }
if (b < 0) {
                                                               5.16 NTT
 11u b2 = (b \% m + m) \% m;
                                                               template <int mod, int G, int maxn>
 ans -= 1ULL * n * ((b2 - b) / m);
                                                               struct NTT {
 b = b2:
                                                                static_assert (maxn == (maxn & -maxn));
                                                                int roots[maxn];
return ans + floor_sum_unsigned(n, m, a, b);
                                                                NTT () {
                                                                 int r = modpow(G, (mod - 1) / maxn);
                                                                 for (int i = maxn >> 1; i; i >>= 1) {
5.13 FWT
                                                                  roots[i] = 1;
/* xor convolution:
                                                                  for (int j = 1; j < i; j++)
* x = (x0,x1) , y = (y0,y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                   roots[i + j] = modmul(roots[i + j - 1], r);
                                                                   r = modmul(r, r);
* x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
* z = (1/2) * z''
                                                                // n must be 2^k, and 0 \le F[i] < mod
                                                                void operator()(int F[], int n, bool inv = false) {
* or convolution:
                                                                 for (int i = 0, j = 0; i < n; i++) {
  if (i < j) swap(F[i], F[j]);</pre>
* x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                                  for (int k = n > 1; (j^k < k) < k; k > = 1);
 * and convolution:
* x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
const LL MOD = 1e9+7;
                                                                 for (int s = 1; s < n; s *= 2) {
                                                                  for (int i = 0; i < n; i += s * 2) {
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
for( int d = 1 ; d < N ; d <<= 1 ) {</pre>
                                                                    for (int j = 0; j < s; j++) {
  int d2 = d << 1;
                                                                    int a = F[i+j];
  for( int s = 0 ; s < N ; s += d2 )
                                                                    int b = modmul(F[i+j+s], roots[s+j]);
  for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
  LL ta = x[ i ] , tb = x[ j ];
  x[ i ] = ta+tb;</pre>
                                                                    F[i+j] = modadd(a, b); // a + b
                                                                    F[i+j+s] = modsub(a, b); // a - b
   x[ j ] = ta-tb;
   if( x[ i ] >= MOD ) x[ i ] -= MOD;
    if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
                                                                 if (inv) {
   }
                                                                  int invn = modinv(n);
                                                                  for (int i = 0; i < n; i++)
if( inv )
                                                                   F[i] = modmul(F[i], invn);
                                                                   reverse(F + 1, F + n);
 for( int i = 0 ; i < N ; i++ ) {</pre>
  x[ i ] *= inv( N, MOD );
  x[ i ] %= MOD;
  }
                                                               NTT<2013265921, 31, 1048576> ntt;
5.14
      Gauss Elimination
                                                               5.17 Partition Number
void gauss(vector<vector<double>> &d) {
                                                               int b = sqrt(n);
                                                               ans[0] = tmp[0] = 1;
int n = d.size(), m = d[0].size();
for (int i = 0; i < m; ++i) {
                                                               for (int i = 1; i <= b; i++) {
 int p = -1;
                                                                for (int rep = 0; rep < 2; rep++)</pre>
 for (int j = i; j < n; ++j) {
                                                                 for (int j = i; j <= n - i * i; j++)
  if (fabs(d[j][i]) < eps) continue;</pre>
                                                                  modadd(tmp[j], tmp[j-i]);
                                                                for (int j = i * i; j <= n; j++)
modadd(ans[j], tmp[j - i * i]);</pre>
   if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
 if (p == -1) continue;
 for (int j = 0; j < m'; ++j) swap(d[p][j], d[i][j]); for (int j = 0; j < n; ++j) {
                                                               5.18 Pi Count (Linear Sieve)
  if (i == j) continue;
                                                               static constexpr int N = 1000000 + 5;
   double z = d[j][i] / d[i][i];
                                                               1ld pi[N];
   for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
                                                               vector<int> primes;
                                                               bool sieved[N];
                                                               1ld cube_root(lld x){
                                                                lld s=cbrt(x-static_cast<long double>(0.1));
                                                                while(s*s*s <= x) ++s;
5.15
      Miller Rabin
                                                                return s-1;
bool isprime(llu x)
static auto witn = [](llu a, llu u, llu n, int t) {
                                                               11d square_root(11d x){
 if (!a) return false;
                                                                lld s=sqrt(x-static_cast<long double>(0.1));
                                                                while(s*s \ll x) ++s;
 while (t--) {
  11u a2 = mmul(a, a, n);
                                                                return s-1;
  if (a2 == 1 && a != 1 && a != n - 1) return true;
  a = a2;
                                                               void init(){
                                                                primes.reserve(N);
 return a != 1;
                                                                primes.push_back(1);
                                                                for(int i=2;i<N;i++) {</pre>
if (x < 2) return false;
                                                                 if(!sieved[i]) primes.push_back(i);
if (!(x & 1)) return x == 2;
                                                                 pi[i] = !sieved[i] + pi[i-1];
 int t = __builtin_ctzll(x - 1);
                                                                 for(int p: primes) if(p > 1) {
                                                                  if(p * i >= N) break;
```

```
sieved[p * i] = true;
                                                                    if (size() == 1) return V{modinv(*begin())};
   if(p % i == 0) break;
                                                                    const int sz = n2k(size() * 2);
                                                                    Poly X = Poly(*this, (size() + 1) / 2).Inv().isz(sz),
                                                                    Y(*this, sz);
ntt(X.data(), sz), ntt(Y.data(), sz);
                                                                    fi(0, sz) X[i] = modmul(X[i], modsub(2, modmul(X[i],
11d phi(11d m, 11d n) {
 static constexpr int MM = 80000, NN = 500;
                                                                      Y[i])));
 static 1ld val[MM][NN];
                                                                    ntt(X.data(), sz, true);
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
                                                                    return X.isz(size());
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
                                                                   Poly Sqrt() const { // coef[0] \in [1, mod)^2
 lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
                                                                    if (size() == 1) return V{QuadraticResidue((*this)
 if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
                                                                      [0], mod)};
 return ret;
                                                                    Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
                                                                      size());
                                                                    return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
1ld pi_count(1ld);
11d P2(11d m, 11d n) {
                                                                       + 1);
11d sm = square_root(m), ret = 0;
 for(lld i = n+1;primes[i]<=sm;i++)</pre>
                                                                   pair<Poly, Poly> DivMod(const Poly &rhs) const {
  ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
                                                                    if (size() < rhs.size()) return {V{0}, *this};</pre>
                                                                    const int sz = size() - rhs.size() + 1;
 return ret;
                                                                    Poly X(rhs); X.irev().isz(sz);
11d pi_count(11d m) {
                                                                    Poly Y(*this); Y.irev().isz(sz);
if(m < N) return pi[m];</pre>
                                                                    Poly Q = Y.Mul(X.Inv()).isz(sz).irev();
                                                                    X = rhs.Mul(Q), Y = *this;
fi(0, size()) Y[i] = modsub(Y[i], X[i]);
 11d n = pi_count(cube_root(m));
 return phi(m, n) + n - 1 - P2(m, n);
                                                                    return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
5.19 Pollard Rho
                                                                   Poly Dx() const {
  Poly ret(size() - 1);
// does not work when n is prime
                                                                    fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
// return any non-trivial factor
llu pollard_rho(llu n) {
                                                                      1]);
 static auto f = [](llu x, llu k, llu m) {
                                                                    return ret.isz(max<int>(1, ret.size()));
    return add(k, mul(x, x, m), m); };
 if (!(n & 1)) return 2;
                                                                   Poly Sx() const {
 mt19937 rnd(120821011);
                                                                    Poly ret(size() + 1);
                                                                    fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
 while (true) {
  llu y = 2, yy = y, x = rnd() % n, t = 1;
                                                                      this)[i]);
  for (llu sz = 2; t == 1; sz <<= 1, y = yy) {
  for (llu i = 0; t == 1 && i < sz; ++i) {
                                                                    return ret:
    yy = f(yy, x, n);
                                                                   Poly Ln() const { // coef[0] == 1
                                                                    return Dx().Mul(Inv()).Sx().isz(size());
    t = gcd(yy > y ? yy - y : y - yy, n);
                                                                   Poly Exp() const \{ // coef[0] == 0 \}
                                                                    if (size() == 1) return V{1};
  if (t != 1 && t != n) return t;
                                                                    Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size)
                                                                      ());
                                                                    Poly Y = X.Ln(); Y[0] = mod - 1;
fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
5.20
      Polynomial Operations
                                                                    return X.Mul(Y).isz(size());
using V = vector<int>;
#define fi(1, r) for (int i = int(1); i < int(r); ++i)</pre>
template <int mod, int G, int maxn> struct Poly : V {
                                                                   Poly Pow(const string &K) const {
                                                                    int nz = 0;
static uint32_t n2k(uint32_t n) {
  if (n <= 1) return 1;
                                                                    while (nz < size() && !(*this)[nz]) ++nz;</pre>
                                                                    int nk = 0, nk2 = 0;
  return 1u << (32 - __builtin_clz(n - 1));</pre>
                                                                    for (char c : K) {
                                                                     nk = (nk * 10 + c - '0') % mod;
 static NTT<mod,G,maxn> ntt; // coefficients in [0, P)
                                                                     nk2 = nk2 * 10 + c - '0';
 explicit Poly(int n = 1) : V(n) {}
                                                                     if (nk2 * nz >= size())
 Poly(const V &v) : V(v) {}
 Poly(const Poly &p, size_t n) : V(n) {
                                                                      return Poly(size());
                                                                     nk2 %= mod - 1;
  copy_n(p.data(), min(p.size(), n), data());
                                                                    if (!nk && !nk2) return Poly(V{1}, size());
 Poly &irev() { return reverse(data(), data() + size())
      *this; }
                                                                    Poly X = V(data() + nz, data() + size() - nz * (nk2 - nz)
 Poly &isz(int sz) { return resize(sz), *this; }
                                                                       1));
 Poly &iadd(const Poly &rhs) { // n() == rhs.n()
                                                                    int x0 = X[0]:
                                                                    return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
  fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
                                                                      modpow(x0, nk2)).irev().isz(size()).irev();
  return *this;
                                                                   Poly InvMod(int L) { // (to evaluate linear recursion)
Poly R{1, \theta}; // *this * R mod x^L = 1 (*this[\theta] ==
 Poly &imul(int k) {
 fi(0, size())(*this)[i] = modmul((*this)[i], k);
  return *this;
                                                                    for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                                     Poly 0 = R.Mul(Poly(data(), min<int>(2 << level,
 Poly Mul(const Poly &rhs) const {
                                                                      size())));
  const int sz = n2k(size() + rhs.size() - 1);
  Poly X(*this, sz), Y(rhs, sz);
                                                                     Poly Q(2 << level); Q[0] = 1;
   \begin{array}{ll} \mathsf{ntt}(\mathsf{X}.\mathsf{data}(),\;\mathsf{sz}),\;\mathsf{ntt}(\mathsf{Y}.\mathsf{data}(),\;\mathsf{sz}) \\ \mathsf{fi}(\mathsf{0},\;\mathsf{sz})\;\mathsf{X}[\mathtt{i}] \; = \; \mathsf{modmul}(\mathsf{X}[\mathtt{i}],\;\mathsf{Y}[\mathtt{i}]); \end{array} 
                                                                     for (int j = (1 << level); j < (2 << level); ++j)</pre>
                                       sz):
                                                                      Q[j] = modsub(mod, O[j]);
  ntt(X.data(), sz, true);
                                                                     R = R.Mul(Q).isz(4 << level);
  return X.isz(size() + rhs.size() - 1);
                                                                    return R.isz(L);
 Poly Inv() const { // coef[0] != 0
```

if (d[i][s] < eps) continue;</pre>

```
static int LinearRecursion(const V &a, const V &c,
                                                                         if (r == -1 ||
                                                                          d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
    int64_t n) { // a_n = \sum_{i=1}^{n} a_i(n-j)}
  const int k = (int)a.size();
  assert((int)c.size() == k + 1);
Poly C(k + 1), W({1}, k), M = {0, 1};
                                                                        if (r == -1) return false;
                                                                        pivot(r, s);
  fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
  C[k] = 1;
                                                                     VD solve(const VVD &a, const VD &b, const VD &c) {
  while (n) {
   if (n % 2) W = W.Mul(M).DivMod(C).second;
                                                                      m = b.size(), n = c.size();
d = VVD(m + 2, VD(n + 2));
   n /= 2, M = M.Mul(M).DivMod(C).second;
                                                                       for (int i = 0; i < m; ++i)
                                                                       for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
  fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
                                                                       p.resize(m), q.resize(n + 1);
  return ret;
                                                                       for (int i = 0; i < m; ++i)
                                                                       p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
                                                                       for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
#undef fi
                                                                       q[n] = -1, d[m + 1][n] = 1;
using Poly_t = Poly<998244353, 3, 1 << 20>;
                                                                       int r = 0;
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
                                                                       for (int i = 1; i < m; ++i)
                                                                        if (d[i][n + 1] < d[r][n + 1]) r = i;
5.21 Quadratic residue
                                                                       if (d[r][n + 1] < -eps) {</pre>
struct S {
                                                                        pivot(r, n);
 int MOD, w;
                                                                        if (!phase(1) || d[m + 1][n + 1] < -eps)
 int64_t x, y;
                                                                         return VD(n, -inf);
                                                                        for (int i = 0; i < m; ++i) if (p[i] == -1) {
 S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
 : MOD(m), w(w_), x(x_), y(y_) {}
S operator*(const S &rhs) const {
                                                                         int s = min_element(d[i].begin(), d[i].end() - 1)
                                                                              - d[i].begin();
  int w_{-} = w;
                                                                         pivot(i, s);
  if (w_ == -1) w_ = rhs.w;
  assert(w_! = -1 \text{ and } w_ == rhs.w);
  return { MOD, w_,
(x * rhs.x + y * rhs.y % MOD * w) % MOD,
                                                                       if (!phase(0)) return VD(n, inf);
                                                                       VD x(n);
   (x * rhs.y + y * rhs.x) % MOD };
                                                                       for (int i = 0; i < m; ++i)
                                                                       if (p[i] < n) \times [p[i]] = d[i][n + 1];
                                                                       return x;
                                                                     }}
int get_root(int n, int P) {
  if (P == 2 \text{ or } n == 0) \text{ return } n;
                                                                      5.23
                                                                             Simplex Construction
  if (qpow(n, (P - 1) / 2, P) != 1) return -1;
  auto check = [&](int x) {
                                                                     Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that for all 1 \leq j \leq m,
    return qpow(x, (P - 1) / 2, P); };
                                                                     \sum_{1 \le i \le n} A_{ji} x_i \le b_j and x_i \ge 0 for all 1 \le i \le n.
  if (check(n) == P-1) return -1
  int64_t a; int w; mt19937 rnd(7122);
                                                                        1. In case of minimization, let c_i' = -c_i
  do { a = rnd() % P;
w = ((a * a - n) % P + P) % P;
                                                                        2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
  } while (check(w) != P - 1);
  return qpow(S(P, w, a, 1), (P + 1) / 2).x;
                                                                        3. \sum_{1 < i < n} A_{ji} x_i = b_j
                                                                              • \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
5.22 Simplex
                                                                              • \sum_{1 \le i \le n} A_{ji} x_i \ge b_j
namespace simplex {
// maximize c^Tx under Ax <= B
                                                                        4. If x_i has no lower bound, replace x_i with x_i - x_i'
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
                                                                           Geometry
using VD = vector<double>;
using VVD = vector<vector<double>>;
                                                                     6.1 Basic Geometry
const double eps = 1e-9;
const double inf = 1e+9;
                                                                     #define IM imag
int n, m;
                                                                     #define RE real
VVD d;
                                                                     using lld = int64_t;
vector<int> p, q;
                                                                     using llf = long double;
void pivot(int r, int s) {
  double inv = 1.0 / d[r][s];
                                                                     using PT = std::complex<lld>;
                                                                     using PTF = std::complex<llf>
                                                                     auto toPTF(PT p) { return PTF{RE(p), IM(p)}; }
int sgn(lld x) { return (x > 0) - (x < 0); }
lld dot(PT a, PT b) { return RE(conj(a) * b); }</pre>
 for (int i = 0; i < m + 2; ++i)
  for (int j = 0; j < n + 2; ++j)
if (i != r && j != s)
    d[i][j] -= d[r][j] * d[i][s] * inv;
                                                                     11d cross(PT a, PT b) { return IM(conj(a) * b); }
for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;
d[r][s] = inv; swap(p[r], q[s]);</pre>
                                                                     int ori(PT a, PT b, PT c) {
                                                                       return sgn(cross(b - a, c - a));
                                                                     bool operator<(const PT &a, const PT &b) ·
                                                                      return RE(a) != RE(b) ? RE(a) < RE(b) : IM(a) < IM(b);</pre>
bool phase(int z) {
 int x = m + z
                                                                     int quad(PT p) {
 while (true) {
                                                                      return (IM(p) == 0) // use sgn for PTF
  int s = -1;
  for (int i = 0; i <= n; ++i) {
                                                                        ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
   if (!z && q[i] == -1) continue;
   if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
                                                                     int argCmp(PT a, PT b) {
                                                                       // -1 / 0 / 1 <-> < / == / > (atan2)
  if (d[x][s] > -eps) return true;
                                                                       int qa = quad(a), qb = quad(b);
  int r = -1;
for (int i = 0; i < m; ++i) {
                                                                      if (qa != qb) return sgn(qa - qb);
```

return sgn(cross(b, a));

};

```
template <typename V> llf area(const V & pt) {
                                                            Point ver(Point a, Point b, Point c) {
                                                             return (b - a) * (c - a);
11d ret = 0;
for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
                                                            vector<Face> convex_hull_3D(const vector<Point> pt) {
 ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
                                                             int n = SZ(pt), ftop = 0
 return ret / 2.0;
                                                             REP(i,n) REP(j,n) flag[i][j] = 0;
                                                             vector<Face> now;
PT rot90(PT p) { return PT{-IM(p), RE(p)}; }
                                                             now.emplace_back(0,1,2);
PTF project(PTF p, PTF q) { // p onto q
                                                             now.emplace_back(2,1,0);
return dot(p, q) * q / dot(q, q);
                                                             for (int i=3; i<n; i++){
                                                              ftop++; vector<Face> next;
REP(j, SZ(now)) {
11f FMOD(11f x) {
if (x < -PI) x += PI * 2;
                                                               Face& f=now[j]; int ff = 0;
if (x > PI) x -= PI * 2;
                                                               ld d=(pt[i]-pt[f.a]).dot(
return x;
                                                                 ver(pt[f.a], pt[f.b], pt[f.c]));
                                                               if (d <= 0) next.push_back(f);</pre>
                                                               if (d > 0) ff=ftop;
6.2 Segment & Line Intersection
                                                               else if (d < 0) ff=-ftop
                                                               flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
struct Segment {
PT st, dir; // represent st + t*dir for 0<=t<=1
                                                              REP(j, SZ(now)) {
Segment(PT s, PT e) : st(s), dir(e - s) {}
                                                               Face& f=now[j]
static bool valid(lld p, lld q) {
                                                               if (flag[f.a][f.b] > 0 &&
  // is there t s.t. 0 <= t <= 1 && qt == p ?
                                                                 flag[f.a][f.b] != flag[f.b][f.a])
 if (q < 0) q = -q, p = -p;
                                                                next.emplace_back(f.a,f.b,i);
  return 0 <= p && p <= q;
                                                               if (flag[f.b][f.c] > 0 &&
                                                                 flag[f.b][f.c] != flag[f.c][f.b])
                                                                next.emplace_back(f.b,f.c,i);
bool isInter(Segment A, PT P) {
                                                               if (flag[f.c][f.a] > 0 &&
if (A.dir == PT(0)) return P == A.st;
                                                                 flag[f.c][f.a] != flag[f.a][f.c])
return cross(P - A.st, A.dir) == 0 &&
                                                                next.emplace_back(f.c,f.a,i);
 Segment::valid(dot(P - A.st, A.dir), norm(A.dir));
                                                              now=next:
template <typename U, typename V>
bool isInter(U A, V B) {
  if (cross(A.dir, B.dir) == 0)
                                                             return now;
  return // handle parallel yourself
   isInter(A, B.st) || isInter(A, B.st+B.dir) ||
                                                            6.5 2D Farthest Pair
isInter(B, A.st) || isInter(B, A.st+A.dir);
PT D = B.st - A.st;
                                                            // stk is from convex hull
11d C = cross(A.dir, B.dir);
                                                            n = (int)(stk.size());
 return U::valid(cross(D, A.dir), C) &&
                                                            int pos = 1, ans = 0; stk.push_back(stk[0]);
   V::valid(cross(D, B.dir), C);
                                                            for(int i=0;i<n;i++)</pre>
                                                             for(int i=0;i<n;i++) {
while(abs(cross(stk[i+1]-stk[i],
struct Line
                                                               stk[(pos+1)%n]-stk[i])) >
PT st, ed, dir;
                                                               abs(cross(stk[i+1]-stk[i]
Line (PT s, PT e)
                                                               stk[pos]-stk[i]))) pos = (pos+1)%n;
  : st(s), ed(e), dir(e - s) {}
                                                             ans = max({ans, dis(stk[i], stk[pos]),
                                                              dis(stk[i+1], stk[pos])});
PTF intersect(const Line &A, const Line &B) {
11f t = cross(B.st - A.st, B.dir) /
                                                            6.6 2D Closest Pair
 llf(cross(A.dir, B.dir));
return toPTF(A.st) + PTF(t) * toPTF(A.dir);
                                                            struct cmp_y {
                                                             bool operator()(const P& p, const P& q) const {
                                                              return p.y < q.y;</pre>
6.3 2D Convex Hull
// returns a convex hull in counterclockwise order
// for a non-strict one, change cross >= to >
                                                            multiset<P, cmp_y> s;
vector<PT> convex_hull(vector<PT> p) {
                                                            void solve(P a[], int n) {
                                                             sort(a, a + n, [](const P& p, const P& q) {
sort(all(p));
 if (p[0] == p.back()) return {p[0]};
                                                              return tie(p.x, p.y) < tie(q.x, q.y);</pre>
int n = p.size(), t = 0;
vector<PT> h(n + 1);
                                                             llf d = INF; int pt = 0;
 for (int _ = 2, s = 0; _--; s = --t, reverse(all(p)))
                                                             for (int i = 0; i < n; ++i) {
 for (PT i : p) {
                                                              while (pt < i and a[i].x - a[pt].x >= d)
                                                               s.erase(s.find(a[pt++]));
   while (t > s + 1 \&\& cross(i, h[t-1], h[t-2]) >= 0)
                                                              auto it = s.lower_bound(P(a[i].x, a[i].y - d));
                                                              while (it != s.end() and it->y - a[i].y < d)
  h[t++] = i;
                                                               d = min(d, dis(*(it++), a[i]));
return h.resize(t), h;
                                                              s.insert(a[i]);
6.4 3D Convex Hull
                                                                  kD Closest Pair (3D ver.)
// return the faces with pt indexes
int flag[MXN][MXN];
                                                            11f solve(vector<P> v) {
                                                             shuffle(v.begin(), v.end(), mt19937());
struct Point{
                                                             unordered_map<lld, unordered_map<lld,
ld x, y, z;
Point operator * (const ld &b) const {
                                                              unordered_map<lld, int>>> m;
                                                             llf d = dis(v[0], v[1]);
  return (Point) {x*b, y*b, z*b};}
Point operator * (const Point &b) const {
                                                             auto Idx = [&d] (11f x) -> 11d {
                                                              return round(x * 2 / d) + 0.1; };
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
                                                             auto rebuild_m = [&m, &v, &Idx](int k) {
```

m.clear();

if (que.size() <= 1 ||</pre>

```
for (int i = 0; i < k; ++i)
                                                                   argCmp(que.front().dir, que.back().dir) == 0)
   m[Idx(v[i].x)][Idx(v[i].y)]
                                                                   return 0:
                                                                 pt.push_back(intersect(que.front(), que.back()));
    [Idx(v[i].z)] = i;
 }; rebuild_m(2);
                                                                 return area(pt);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
                                                              6.10 Minkowski Sum
     kz = Idx(v[i].z); bool found = false;
  for (int dx = -2; dx <= 2; ++dx) {
                                                              vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
   const 11d nx = dx + kx
                                                                hull(A), hull(B);
   if (m.find(nx) == m.end()) continue;
                                                               vector<pll> C(1, A[0] + B[0]), s1, s2;
for(int i = 0; i < SZ(A); ++i)
s1.pb(A[(i + 1) % SZ(A)] - A[i]);</pre>
   auto& mm = m[nx];
   for (int dy = -2; dy <= 2; ++dy) {
    const 11d ny = dy + ky;
                                                                for(int i = 0; i < SZ(B); i++)</pre>
    if (mm.find(ny) == mm.end()) continue;
                                                                 s2.pb(B[(i + 1) % SZ(B)] - B[i]);
    auto& mmm = mm[ny];
                                                                for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
    for (int dz = -2; dz <= 2; ++dz) {
                                                                 if (p2 >= SZ(B)
     const 11d nz = dz + kz;
                                                                   || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
     if (mmm.find(nz) == mmm.end()) continue;
                                                                  C.pb(C.back() + s1[p1++]);
     const int p = mmm[nz];
     if (dis(v[p], v[i]) < d) {
  d = dis(v[p], v[i]);</pre>
                                                                  C.pb(C.back() + s2[p2++]);
                                                                return hull(C), C;
      found = true;
                                                              6.11 Circle Class
                                                              struct Circle { PTF o; llf r; };
  if (found) rebuild_m(i + 1);
                                                              vector<llf> intersectAngle(Circle A, Circle B) {
  else m[kx][ky][kz] = i;
                                                               PTF dir = B.o - A.o; 11f d2 = norm(dir);
                                                               if (norm(A.r - B.r) >= d2) // norm(x) := |x|^2
 return d;
                                                                if (A.r < B.r) return {-PI, PI}; // A in B</pre>
                                                                 else return {}; // B in A
                                                                if (norm(A.r + B.r) <= d2) return {};</pre>
6.8 Simulated Annealing
                                                                llf dis = abs(dir), theta = arg(dir);
11f anneal() {
                                                                11f phi = acos((A.r * A.r + d2 - B.r * B.r) /
 mt19937 rnd_engine( seed );
                                                                  (2 * A.r * dis));
 uniform_real_distribution< llf > rnd( 0, 1 );
                                                               11f L = FMOD(theta - phi), R = FMOD(theta + phi);
 const llf dT = 0.001;
                                                               return { L, R };
 // Argument p
11f S_cur = calc( p ), S_best = S_cur;
for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
                                                               vector<PTF> intersectPoint(Circle a, Circle b) {
 // Modify p to p_prime
const llf S_prime = calc( p_prime );
                                                               llf d = abs(a.o - b.o);
                                                                if (d >= b.r+a.r || d <= abs(b.r-a.r)) return {};</pre>
  const llf delta_c = S_prime - S_cur
                                                                11f dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
  11f prob = min( ( 11f ) 1, exp( -delta_c / T ) );
                                                               PTF dir = (a.o - b.o) / d;
  if ( rnd( rnd_engine ) <= prob )</pre>
                                                                PTF u = dir*d1 + b.o;
   S_cur = S_prime, p = p_prime;
                                                               PTF v = rot90(dir) * sqrt(max<llf>(0, b.r*b.r-d1*d1));
  if ( S_prime < S_best ) // find min</pre>
                                                                return \{u + v, u - v\};
   S_best = S_prime, p_best = p_prime;
return S_best;
                                                              6.12 Intersection of line and Circle
                                                              vector<PTF> line_interCircle(const PTF &p1,
6.9 Half Plane Intersection
                                                                 const PTF &p2, const PTF &c, const double r)
                                                                PTF ft = p1 + project(c-p1, p2-p1), vec = p2-p1;
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
                                                                llf dis = abs(c - ft);
bool operator<(const Line &lhs, const Line &rhs) {</pre>
                                                                if (abs(dis - r) < eps) return {ft};</pre>
  if (int cmp = argCmp(lhs.dir, rhs.dir))
                                                                if (dis > r) return {};
    return cmp == -1;
                                                                vec = vec * sqrt(r * r - dis * dis) / abs(vec);
  return ori(lhs.st, lhs.ed, rhs.st) < 0;</pre>
                                                                return {ft + vec, ft - vec};
// intersect function is in "Segment Intersect"
                                                              6.13 Intersection of Polygon and Circle
11f HPI(vector<Line> &lines) {
  sort(lines.begin(), lines.end());
                                                              // Divides into multiple triangle, and sum up
                                                              // test by HDU2892
  deque<Line> que;
  deque<PTF> pt;
                                                              11f _area(PTF pa, PTF pb, llf r)
  que.push_back(lines[0]);
                                                               if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
                                                               if (abs(pb) < eps) return 0;
11f S, h, theta;</pre>
  for (int i = 1; i < (int)lines.size(); i++) {</pre>
    if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
                                                                11f a = abs(pb), b = abs(pa), c = abs(pb - pa);
     continue
#define POP(L, R) \
                                                                11f cosB = dot(pb, pb - pa) / a / c, B = acos(cosB);
                                                                11f \cos C = dot(pa, pb) / a / b, C = acos(cosC);
    while (pt.size() > 0 \
      && ori(L.st, L.ed, pt.back()) < 0) \
                                                                if (a > r)
    pt.pop_back(), que.pop_back(); \
while (pt.size() > 0 \
                                                                 S = (C / 2) * r * r;
                                                                 h = a * b * sin(C) / c;
      && ori(R.st, R.ed, pt.front()) < 0) \
                                                                 if (h < r && B < PI / 2)
      pt.pop_front(), que.pop_front();
                                                                  S = (acos(h / r) * r * r - h * sqrt(r*r - h*h));
                                                               } else if (b > r) {
  theta = PI - B - asin(sin(B) / r * a);
    POP(lines[i], lines[i])
    pt.push_back(intersect(que.back(), lines[i]));
    que.push_back(lines[i]);
                                                                 S = 0.5 * a*r * sin(theta) + (C - theta) / 2 * r*r;
                                                                } else
  POP(que.front(), que.back())
                                                                 S = 0.5 * sin(C) * a * b;
```

return S;

```
11f area_poly_circle(const vector<PTF> &poly,
 const PTF &0, const llf r) {
11f S = 0:
for (int i = 0, N = poly.size(); i < N; ++i)</pre>
 S += _area(poly[i] - 0, poly[(i + 1) % N] - 0, r) *
  ori(0, poly[i], poly[(i + 1) % N]);
return fabs(S);
6.14 Point & Hulls Tangent
#define above(P, Vi, Vj) (ori(P, Vi, Vj) > 0) // true
    if Vi is above Vj
#define below(P, Vi, Vj) (ori(P, Vi, Vj) < 0) // true</pre>
    if Vi is below Vj
// Rtangent_PointPolyC(): binary search for convex
    polygon right tangent
   Input: P = a 2D point (exterior to the polygon)
        n = number of polygon vertices
//
        V = array of vertices for a 2D convex polygon
    with V[n] = V[0]
// Return: index "i" of rightmost tangent point V[i]
int Rtangent_PointPolyC(PT P, int n, PT *V) {
int a, b, c;
int upA, dnC;
if (below(P, V[1], V[0]) && !above(P, V[n - 1], V[0]))
 return 0;
for (a = 0, b = n;;) {
 c = (a + b) / 2;
 dnC = below(P, V[c + 1], V[c]);
if (dnC && !above(P, V[c - 1], V[c]))
  return c:
 upA = above(P, V[a + 1], V[a]);
  if (upA) {
  if (dnC) {
   b = c;
   } else {
    if (above(P, V[a], V[c]))
     b = c;
    else
     a = c;
   }
  } else {
   if (!dnC) {
   a = c;
   } else {
    if (below(P, V[a], V[c]))
    b = c;
    else
     a = c:
   }
// Ltangent_PointPolyC(): binary search for convex
    polygon left tangent
    Input: P = a 2D point (exterior to the polygon)
//
        n = number of polygon vertices
//
        V = array of vertices for a 2D convex polygon
    with V[n]=V[0]
   Return: index "i" of leftmost tangent point V[i]
int Ltangent_PointPolyC(PT P, int n, PT *V) {
int a, b, c;
int dnA, dnC;
if (above(P, V[n - 1], V[0]) && !below(P, V[1], V[0]))
 return 0;
for (a = 0, b = n;;) {
 c = (a + b) / 2
 dnC = below(P, V[c + 1], V[c]);
  if (above(P, V[c - 1], V[c]) && !dnC)
   return c;
  dnA = below(P, V[a + 1], V[a]);
  if (dnA) {
  if (!dnC) {
   b = c;
   } else {
```

```
if (below(P, V[a], V[c]))
   b = c;
else
   a = c;
}
} else {
   if (dnC) {
    a = c;
} else {
   if (above(P, V[a], V[c]))
    b = c;
else
   a = c;
}
}
}
```

6.15 Convex Hulls Tangent

```
// RLtangent_PolyPolyC(): get the RL tangent between
    two convex polygons
    Input: m = number of vertices in polygon 1
//
        V = array of vertices for convex polygon 1 with
     V[m]=V[0]
        n = number of vertices in polygon 2
//
        W = array of vertices for convex polygon 2 with
     W[n]=W[0]
    Output: *t1 = index of tangent point V[t1] for
//
    polygon 1
11
        *t2 = index of tangent point W[t2] for polygon
void RLtangent_PolyPolyC(int m, PT *V, int n, PT *W,
    int *t1, int *t2) {
 int ix1, ix2; // search indices for polygons 1 and 2
 // first get the initial vertex on each polygon
 ix1 = Rtangent_PointPolyC(W[0], m, V); // right
    tangent from W[0] to V
 ix2 = Ltangent_PointPolyC(V[ix1], n, W); // left
    tangent from V[ix1] to W
 // ping-pong linear search until it stabilizes
 int done = false; // flag when done
 while (done == false) {
  done = true; // assume done until..
  while (ori(W[ix2], V[ix1], V[ix1 + 1]) <= 0) {</pre>
   ++ix1; // get Rtangent from W[ix2] to V
  while (ori(V[ix1], W[ix2], W[ix2 - 1]) >= 0) {
             // get Ltangent from V[ix1] to W
   --ix2:
   done = false; // not done if had to adjust this
  }
 *t1 = ix1;
 *t2 = ix2;
 return;
```

6.16 Tangent line of Two Circle

```
vector<Line>
tanline(const Circle &c1, const Circle &c2, int sign1){
 // sign1 = 1 for outer tang, -1 for inter tang
 vector<Line> ret;
 if (norm(c1.o - c2.o) < eps) return ret;</pre>
 11f d = abs(c1.o - c2.o);
 PTF v = (c2.o - c1.o) / d;
llf c = (c1.r - sign1 * c2.r) / d;
 if (c * c > 1) return ret;
 llf h = sqrt(max<llf>(0, 1 - c * c));
 for (int sign2 : {1, -1}) {
  PTF n = c * v + sign2 * h * rot90(v);
  PTF p1 = c1.o + n * c1.r;
  PTF p2 = c2.o + n * (c2.r * sign1);
  if (norm(p2 - p1) < eps)
   p2 = p1 + rot90(c2.o - c1.o);
  ret.push_back({p1, p2});
 return ret;
```

6.17 Minimum Covering Circle

int touch(Node* r, int x, int y, LL d2){

```
LL dis = sqrt(d2)+1;
                                                                  if (x<r->x1-dis || x>r->x2+dis ||
template<typename P>
                                                                    y<r->y1-dis || y>r->y2+dis)
Circle getCircum(const P &a, const P &b, const P &c){
                                                                    return 0;
Real a1 = a.x-b.x, b1 = a.y-b.y;
                                                                  return 1;
Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
Real a2 = a.x-c.x, b2 = a.y-c.y;
                                                                 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
                                                                  if (!r || !touch(r, x, y, md2)) return;
Circle cc;
                                                                  LL d2 = dis2(r->x, r->y, x, y);
cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
                                                                  if (d2 < md2 \mid | (d2 == md2 && mID < r->id)) {
cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
                                                                   mID = r->id;
cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
                                                                   md2 = d2;
return cc;
                                                                  }
                                                                  // search order depends on split dim
                                                                  if ((r->f == 0 \&\& x < r->x) ||
template<typename P>
                                                                     (r->f == 1 \&\& y < r->y))
Circle MinCircleCover(const vector<P>& pts){
                                                                   nearest(r->L, x, y, mID, md2);
{\tt random\_shuffle(pts.begin(), pts.end());}
                                                                   nearest(r->R, x, y, mID, md2);
Circle c = { pts[0], 0 }
                                                                  } else {
for(int i=0;i<(int)pts.size();i++){</pre>
                                                                   nearest(r->R, x, y, mID, md2);
 if (dist(pts[i], c.o) <= c.r) continue;</pre>
                                                                   nearest(r->L, x, y, mID, md2);
  c = { pts[i], 0 };
 for (int j = 0; j < i; j++) {
  if(dist(pts[j], c.o) <= c.r) continue;</pre>
                                                                 int query(int x, int y) {
   c.o = (pts[i] + pts[j]) / 2;
                                                                  int id = 1029384756;
  c.r = dist(pts[i], c.o);
for (int k = 0; k < j; k++) {</pre>
                                                                  LL d2 = 102938475612345678LL;
                                                                  nearest(root, x, y, id, d2);
   if (dist(pts[k], c.o) <= c.r) continue;</pre>
                                                                  return id;
    c = getCircum(pts[i], pts[j], pts[k]);
                                                                 }
                                                                } tree;
  }
                                                                       Rotating Sweep Line
return c;
                                                                void rotatingSweepLine(pair<int, int> a[], int n) {
                                                                 vector<pair<int, int>> 1;
                                                                 1.reserve(n * (n - 1) / 2);
6.18
      KDTree (Nearest Point)
                                                                 for (int i = 0; i < n; ++i)
const int MXN = 100005;
                                                                  for (int j = i + 1; j < n; ++j)
struct KDTree {
                                                                   1.emplace_back(i, j);
struct Node {
                                                                 sort(1.begin(), 1.end(), [&a](auto &u, auto &v){
                                                                  1ld udx = a[u.first].first - a[u.second].first;
 int x,y,x1,y1,x2,y2;
 int id,f;
Node *L, *R;
                                                                  11d udy = a[u.first].second - a[u.second].second;
                                                                  lld vdx = a[v.first].first - a[v.second].first;
lld vdy = a[v.first].second - a[v.second].second;
 } tree[MXN], *root;
int n;
                                                                  if (udx == 0 or vdx == 0) return not udx == 0;
LL dis2(int x1, int y1, int x2, int y2) {
  LL dx = x1-x2, dy = y1-y2;
                                                                  int s = sgn(udx * vdx);
                                                                  return udy * vdx * s < vdy * udx * s;
 return dx*dx+dy*dy;
                                                                 vector<int> idx(n), p(n);
 static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
                                                                 iota(idx.begin(), idx.end(), 0);
                                                                 sort(idx.begin(), idx.end(), [&a](int i, int j){
static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
                                                                  return a[i] < a[j]; });
void init(vector<pair<int,int>> ip) {
 n = ip.size();
                                                                 for (int i = 0; i < n; ++i) p[idx[i]] = i;
 for (int i=0; i<n; i++) {</pre>
                                                                 for (auto [i, j]: 1) {
  tree[i].id = i;
                                                                  // do here
                                                                  swap(p[i], p[j]);
idx[p[i]] = i, idx[p[j]] = j;
   tree[i].x = ip[i].first;
   tree[i].y = ip[i].second;
                                                                }
 root = build_tree(0, n-1, 0);
                                                                6.20 Circle Cover
Node* build_tree(int L, int R, int d) {
 if (L>R) return nullptr;
int M = (L+R)/2; tree[M].f = d%2;
                                                                const int N = 1021;
                                                                struct CircleCover {
 nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
                                                                 int C;
 tree[M].x1 = tree[M].x2 = tree[M].x;
                                                                 Cir c[N]
                                                                 bool g[N][N], overlap[N][N];
  tree[M].y1 = tree[M].y2 = tree[M].y;
 tree[M].L = build_tree(L, M-1, d+1);
                                                                 // Area[i] : area covered by at least i circles
  if (tree[M].L) {
                                                                 double Area[ N ];
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
                                                                 void init(int _C){ C = _C;}
   tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
                                                                 struct Teve {
  tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
                                                                  PTF p; double ang; int add;
                                                                  Teve() {}
  tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
                                                                  Teve(PTF _a, double _b, int _c):p(_a), ang(_b), add(
 tree[M].R = build_tree(M+1, R, d+1);
                                                                     _c){}
 if (tree[M].R) {
                                                                  bool operator<(const Teve &a)const
  tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
                                                                  {return ang < a.ang;}
  tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
                                                                 }eve[N * 2];
                                                                 // strict: x = 0, otherwise x = -1
bool disjuct(Cir &a, Cir &b, int x)
  tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
                                                                 {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
                                                                 bool contain(Cir &a, Cir &b, int x)
{return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
  return tree+M;
```

bool contain(int i, int j) {

memcpy(x, c, sizeof(int) * z);

```
/* c[j] is non-strictly in c[i]. */
                                                                   for (int i = n - 1; i >= 0; --i)
if (a[i] && t[a[i] - 1])
  return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c
    [j].R) == 0 \&\& i < j)) \&\& contain(c[i], c[j], -1);
                                                                     a[--x[s[a[i] - 1]]] = a[i] - 1;
                                                                  void_sais(int *s, int *a, int *p, int *q,
 void solve(){
  fill_n(Area, C + 2, 0);
                                                                   bool *t, int *c, int n, int z) {
  for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)</pre>
                                                                   bool uniq = t[n - 1] = true;
                                                                   int nn=0, nmxz=-1, *nsa = a+n, *ns=s+n, last=-1;
                                                                   memset(c, 0, sizeof(int) * z);
    overlap[i][j] = contain(i, j);
  for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)</pre>
                                                                   for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
    g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                                                                   if (uniq) {
       disjuct(c[i], c[j], -1));
                                                                    for (int i = 0; i < n; ++i) a[--c[s[i]]] = i;
  for(int i = 0; i < C; ++i){</pre>
                                                                    return;
   int E = 0, cnt = 1;
   for(int j = 0; j < C; ++j)
                                                                   for (int i = n - 2; i \ge 0; --i)
    if(j != i && overlap[j][i])
                                                                    t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
     ++cnt;
                                                                   pre(a, c, n, z);
   for(int j = 0; j < C; ++j)
                                                                   for (int i = 1; i <= n - 1; ++i)
    if(i != j && g[i][j]) {
                                                                    if (t[i] && !t[i - 1])
     auto IP = intersectPoint(c[i], c[j]);
                                                                     a[--x[s[i]]] = p[q[i] = nn++] = i;
                                                                   induce(a, c, s, t, n, z);
for (int i = 0; i < n; ++i)
     PTF aa = IP[0], bb = IP[1];
     llf A = arg(aa-c[i].0), B = arg(bb-c[i].0);
                                                                    if (a[i] && t[a[i]] && !t[a[i] - 1]) {
     eve[E++] = Teve(bb,B,1), eve[E++]=Teve(aa,A,-1);
     if(B > A) ++cnt;
                                                                    bool neq = last < 0 || \</pre>
                                                                     memcmp(s + a[i], s + last,
(p[q[a[i]] + 1] - a[i]) * sizeof(int));
   if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
                                                                    ns[q[last = a[i]]] = nmxz += neq;
   else{
    sort(eve, eve + E);
    eve[E] = eve[0];
                                                                   sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
    for(int j = 0; j < E; ++j){
                                                                   pre(a, c, n, z);
                                                                   for (int i = nn - 1; i >= 0; --i)
a[--x[s[p[nsa[i]]]]] = p[nsa[i]];
     cnt += eve[j].add;
     Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
     double theta = eve[j + 1].ang - eve[j].ang;
                                                                   induce(a, c, s, t, n, z);
     if (theta < 0) theta += 2. * pi;</pre>
     Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
                                                                  void build(const string &s) {
                                                                   const int n = int(s.size());
                                                                   for (int i = 0; i < n; ++i) _s[i] = s[i];
                                                                   _s[n] = 0; // s shouldn't contain 0
                                                                   sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
for(int i = 0; i < n; ++i) rev[sa[i] = sa[i+1]] = i;</pre>
                                                                   int ind = hi[0] = 0;
                                                                   for (int i = 0; i < n; ++i) {
7
     Stringology
                                                                    if (!rev[i]) {
                                                                     ind = 0;
7.1
     Hash
                                                                     continue;
class Hash {
 private:
                                                                    while (i + ind < n && \</pre>
  static constexpr int P = 127, Q = 1051762951;
                                                                     s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  vector<int> h, p;
                                                                    hi[rev[i]] = ind ? ind-- : 0;
 public:
  void init(const string &s){
   h.assign(s.size()+1, 0); p.resize(s.size()+1);
   for (size_t i = 0; i < s.size(); ++i)</pre>
                                                                  7.3 Suffix Automaton
    h[i + 1] = add(mul(h[i], P), s[i]);
                                                                  struct SuffixAutomaton {
   generate(p.begin(), p.end(),[x=1,y=1,this]()
mutable{y=x;x=mul(x,P);return y;});
                                                                   struct node -
                                                                    int ch[K], len, fail, cnt, indeg;
                                                                    node(int L = 0) : ch{}, len(L), fail(0), cnt(0),
  int query(int 1, int r){ // 1-base (1, r]
  return sub(h[r], mul(h[1], p[r-1]));}
                                                                       indeg(0) {}
                                                                   } st[N];
                                                                   int root, last, tot;
                                                                   void extend(int c)
7.2 Suffix Array
                                                                    int cur = ++tot;
                                                                    st[cur] = node(st[last].len + 1);
namespace sfx {
bool _t[maxn * 2];
                                                                    while (last && !st[last].ch[c]) {
int hi[maxn], rev[maxn];
                                                                      st[last].ch[c] = cur;
int _s[maxn * 2], sa[maxn * 2], _c[maxn * 2];
int x[maxn], _p[maxn], _q[maxn * 2];
// sa[i]: sa[i]-th suffix is the
                                                                       last = st[last].fail;
                                                                    if (!last) {
                                                                       st[cur].fail = root;
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix
                                                                    } else {
// of suffix sa[i] and suffix sa[i - 1].
                                                                       int q = st[last].ch[c];
void pre(int *a, int *c, int n, int z) {
                                                                       if (st[q].len == st[last].len + 1) {
memset(a, 0, sizeof(int) * n);
                                                                         st[cur].fail = q;
 memcpy(x, c, sizeof(int) * z);
                                                                       } else {
                                                                         int clone = ++tot;
                                                                         st[clone] = st[q];
void induce(int *a,int *c,int *s,bool *t,int n,int z){
memcpy(x + 1, c, sizeof(int) * (z - 1));
                                                                         st[clone].len = st[last].len + 1;
 for (int i = 0; i < n; ++i)
                                                                         st[st[cur].fail = st[q].fail = clone].cnt = 0;
  if (a[i] && !t[a[i] - 1])
                                                                         while (last && st[last].ch[c] == q) {
   a[x[s[a[i] - 1]]++] = a[i] - 1;
                                                                           st[last].ch[c] = clone;
```

last = st[last].fail;

for (int i = 1; i < t.length(); ++i) {</pre>

```
z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1)
                                                               while (i - z[i] >= 0 \&\& i + z[i] < t.length()) {
    }
  }
                                                                if(t[i - z[i]] == t[i + z[i]]) ++z[i];
  st[last = cur].cnt += 1;
                                                                else break:
 void init(const char* s) {
                                                               if (i + z[i] > r) r = i + z[i], l = i;
  root = last = tot = 1;
  st[root] = node(0);
                                                              for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
  for (char c; c = *s; ++s) extend(c - 'a');
                                                              return ans;
 int q[N];
                                                             7.7 Lexico Smallest Rotation
 void dp() {
  for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
                                                             string mcp(string s) {
                                                              int n = s.length();
  int head = 0, tail = 0;
                                                              s += s; int i = 0, j = 1;
  for (int i = 1; i <= tot; i++)
                                                              while (i < n && j < n) {</pre>
    if (st[i].indeg == 0) q[tail++] = i;
                                                               int k = 0;
  while (head != tail) {
                                                               while (k < n \&\& s[i + k] == s[j + k]) k++;
    int now = q[head++];
                                                               ((s[i+k] \le s[j+k]) ? j : i) += k + 1;
    if (int f = st[now].fail) {
                                                               j += (i == j);
      st[f].cnt += st[now].cnt;
      if (--st[f].indeg == 0) q[tail++] = f;
                                                              return s.substr(i < n ? i : j, n);</pre>
  }
                                                             7.8 BWT
 int run(const char* s) {
                                                             struct BurrowsWheeler{
  int now = root;
                                                             #define SIGMA 26
  for (char c; c = *s; ++s) {
                                                             #define BASE 'a'
    if (!st[now].ch[c -= 'a']) return 0;
                                                              vector<int> v[ SIGMA ];
    now = st[now].ch[c];
                                                              void BWT(char* ori, char* res){
                                                               // make ori -> ori + ori
  return st[now].cnt;
                                                               // then build suffix array
} SAM;
                                                              void iBWT(char* ori, char* res){
                                                               for( int i = 0 ; i < SIGMA ; i ++ )</pre>
7.4 KMP
                                                                v[ i ].clear();
vector<int> kmp(const string &s) {
                                                               int len = strlen( ori );
 vector<int> f(s.size(), 0);
                                                               for( int i = 0 ; i < len ; i ++ )</pre>
 /* f[i] = length of the longest prefix
                                                                v[`ori[i] - BASE ].push_back( i );
   (excluding s[0:i]) such that it coincides
                                                               vector<int> a;
   with the suffix of s[0:i] of the same length */
                                                               for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
 /* i + 1 - f[i] is the length of the
                                                                for( auto j : v[ i ] ){
  a.push_back( j );
   smallest recurring period of s[0:i] */
 int k = 0;
                                                                 ori[ ptr ++ ] = BASE + i;
 for (int i = 1; i < (int)s.size(); ++i) {</pre>
  while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
                                                               for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
  if (s[i] == s[k]) ++k;
  f[i] = k;
                                                                ptr = a[ ptr ];
 return f;
                                                               res[ len ] = 0;
vector<int> search(const string &s, const string &t) {
                                                             } bwt;
 // return 0-indexed occurrence of t in s
 vector < int > f = kmp(t), r;
                                                             7.9 Palindromic Tree
 for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
                                                             struct palindromic_tree{
   while(k > 0 \&\& (k==(int)t.size() \mid \mid s[i]!=t[k])) 
                                                              struct node{
   k = f[k - 1]
                                                               int next[26],f,len;
  if (s[i] == t[k]) ++k;
                                                               int cnt, num, st, ed;
  if (k == (int)t.size()) r.push_back(i-t.size()+1);
                                                               node(int 1=0):f(0),len(1),cnt(0),num(0) {
                                                                memset(next, 0, sizeof(next)); }
 return res;
                                                              vector<node> st;
                                                              vector<char> s;
7.5 Z value
                                                              int last,n;
vector<int> Zalgo(const string &s) {
                                                              void init(){
 vector<int> z(s.size(), s.size());
                                                               st.clear();s.clear();last=1; n=0;
 for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
                                                               st.push_back(0);st.push_back(-1);
  int j = clamp(r - i, 0, z[i - 1]);
for (; i + j < z[0] and s[i + j] == s[j]; ++j);
                                                               st[0].f=1;s.push_back(-1); }
                                                              int getFail(int x){
  if (i + (z[i] = j) > r) r = i + z[1 = i];
                                                               while(s[n-st[x].len-1]!=s[n])x=st[x].f;
                                                               return x;}
 return z;
                                                              void add(int c){
}
                                                               s.push_back(c-='a'); ++n;
                                                               int cur=getFail(last);
7.6 Manacher
                                                               if(!st[cur].next[c]){
int z[maxn];
                                                                int now=st.size();
                                                                 st.push_back(st[cur].len+2);
int manacher(const string& s) {
 string t = ".";
                                                                st[now].f=st[getFail(st[cur].f)].next[c];
 for(char c: s) t += c, t += '.';
                                                                st[cur].next[c]=now;
 int 1 = 0, r = 0, ans = 0;
                                                                st[now].num=st[st[now].f].num+1;
```

```
last=st[cur].next[c];
++st[last].cnt;}
void dpcnt() {
  for (int i=st.size()-1; i >= 0; i--)
    st[st[i].f].cnt += st[i].cnt;
}
int size() { return st.size()-2;}
} pt;
int main() {
  string s; cin >> s; pt.init();
  for (int i=0; i<SZ(s); i++) {
   int prvsz = pt.size(); pt.add(s[i]);
   if (prvsz != pt.size()) {
    int r = i, l = r - pt.st[pt.last].len + 1;
    // pal @ [l,r]: s.substr(l, r-l+1)
   }
}
return 0;
}</pre>
```

8 Misc

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|{\rm det}(\tilde{L}_{rr})|.$

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\dots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + d_2 + \ldots + d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \le k \le n$.

8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W\subseteq X, |W|\leq |N_G(W)|$

8.1.7 Euler's planar graph formula

$$V - E + F = C + 1$$
, $E \le 3V - 6$ (?)

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Lucas's theorem

```
 \binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p}, \text{ where } m=m_kp^k+m_{k-1}p^{k-1}+\cdots+m_1p+m_0, and n=n_kp^k+n_{k-1}p^{k-1}+\cdots+n_1p+n_0.
```

8.1.10 Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2)$, find maximum $S\in I_1\cap I_2$. For each iteration, build the directed graph and find a shortest path from s to t.

• $s \to x: S \sqcup \{x\} \in I_1$ • $x \to t: S \sqcup \{x\} \in I_2$ • $y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1$ (y is in the unique circuit of $S \sqcup \{x\}$)

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|$. In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

• $x \to y : S \setminus \{y\} \sqcup \{x\} \in I_2$ (y is in the unique circuit of $S \sqcup \{x\}$)

8.2 DP-opt Condition

8.2.1 totally monotone (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

8.2.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

8.3 Convex 1D/1D DP

```
struct segment {
 int i, 1, r;
 segment() {}
 segment(int a, int b, int c): i(a), l(b), r(c) {}
inline lld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
 dp[0] = 0;
 deque < segment > dq; dq.push_back(segment(0, 1, n));
 for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i)
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
   f(i, dq.back().1) < f(dq.back().i, dq.back().1)
    dq.pop_back();
  if (dq.size())
   int d = 1 << 20, c = dq.back().1;
while (d >>= 1) if (c + d <= dq.back().r)</pre>
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.1 = c + 1;
  if (seg.1 <= n) dq.push_back(seg);</pre>
```

8.4 ConvexHull Optimization

```
mutable int64_t a, b, p;
bool operator<(const L &r) const { return a < r.a; }</pre>
bool operator<(int64_t x) const { return p < x; }</pre>
};
struct DynamicHull : multiset<L, less<>> {
static const int64_t kInf = 1e18;
bool Isect(iterator x, iterator y) {
 auto Div = [](int64_t a, int64_t b) {
    return a / b - ((a ^ b) < 0 && a % b);
  if (y == end()) { x->p = kInf; return false; }
 if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
 else x-p = Div(y-b - x-b, x-a - y-a);
 return x->p >= y->p;
void Insert(int64_t a, int64_t b) {
 auto z = insert(\{a, b, 0\}), y = z++, x = y;
 while (Isect(y, z)) z = erase(z);
 if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
 while ((y = x) != begin() \&\& (--x)->p >= y->p)
  Isect(x, erase(y));
int64_t Query(int64_t x) {
 auto 1 = *lower_bound(x);
 return 1.a * x + 1.b;
```

```
8.5
     Josephus Problem
// n people kill m for each turn
int f(int n, int m) {
int s = 0:
for (int i = 2; i <= n; i++)
 s = (s + m) \% i;
 return s;
// died at kth
int kth(int n, int m, int k){
if (m == 1) return n-1;
for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
return k;
8.6 Cactus Matching
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
dfn[u]=low[u]=++dfs_idx;
for(int i=0;i<(int)init_g[u].size();i++){</pre>
 int v=init_g[u][i];
  if(v==par[u]) continue;
  if(!dfn[v]){
  par[v]=u;
   tarjan(v);
  low[u]=min(low[u],low[v]);
  if(dfn[u]<low[v]){</pre>
   g[u].push_back(v);
   g[v].push_back(u);
  }else{
  low[u]=min(low[u],dfn[v]);
   if(dfn[v]<dfn[u]){</pre>
    int temp_v=u;
    bcc_id++;
    while(temp_v!=v){
     g[bcc_id+n].push_back(temp_v);
     g[temp_v].push_back(bcc_id+n);
     temp_v=par[temp_v];
   g[bcc_id+n].push_back(v);
   g[v].push_back(bcc_id+n);
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
int dp[maxn][2],min_dp[2][2],tmp[2][2],tp[2];
void dfs(int u,int fa){
if(u<=n){
 for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
   dfs(v,u)
  memset(tp,0x8f,sizeof tp);
   if(v<=n){
    tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
    tp[1]=max(
     dp[u][0]+dp[v][0]+1
     dp[u][1]+max(dp[v][0],dp[v][1])
  }else{
    tp[0]=dp[u][0]+dp[v][0];
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
   dp[u][0]=tp[0],dp[u][1]=tp[1];
}else{
 for(int i=0;i<(int)g[u].size();i++){</pre>
  int v=g[u][i];
  if(v==fa) continue;
  dfs(v,u);
 min_dp[0][0]=0;
 min_dp[1][1]=1;
min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f3f;
 for(int i=0;i<(int)g[u].size();i++){</pre>
  int v=g[u][i];
```

if(v==fa) continue;

memset(tmp,0x8f,sizeof tmp);

```
tmp[0][0]=max(
    min_dp[0][0]+max(dp[v][0],dp[v][1]),
    min_dp[0][1]+dp[v][0]
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
   tmp[1][0]=max(
    \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
    min_dp[1][1]+dp[v][0]
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
   memcpy(min_dp,tmp,sizeof tmp);
  dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
  dp[u][0]=min_dp[0][0];
int main(){
 int m,a,b;
 scanf("%d%d",&n,&m);
 for(int i=0;i<m;i++){</pre>
  scanf("%d%d",&a,&b);
  init_g[a].push_back(b);
  init_g[b].push_back(a);
 par[1]=-1;
 tarjan(1);
 dfs(1,-1);
 printf("%d\n", max(dp[1][0], dp[1][1]));
 return 0;
     Tree Knapsack
int dp[N][K]; PII obj[N];
vector<int> G[N];
void dfs(int u,
                int mx){
 for(int s: G[u]) {
  if(mx < obj[s].first) continue;</pre>
  for(int i=0;i<=mx-obj[s].FF;i++)</pre>
   dp[s][i] = dp[u][i];
  dfs(s, mx - obj[s].first);
  for(int i=obj[s].FF;i<=mx;i++)</pre>
   dp[u][i] = max(dp[u][i],
    dp[s][i - obj[s].FF] + obj[s].SS);
}
    N Queens Problem
vector< int > solve( int n ) {
 // no solution when n=2, 3
 vector< int > ret:
 if ( n % 6 == 2 ) {
for ( int i = 2 ; i <= n ; i += 2 )
   ret.push_back( i )
  ret.push_back( 3 ); ret.push_back( 1 );
  for ( int i = 7 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 5 );
 } else if ( n % 6 == 3 ) {
for ( int i = 4 ; i <= n ; i += 2 )</pre>
   ret.push_back( i )
  ret.push_back( 2 );
  for ( int i = 5 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 1 ); ret.push_back( 3 );
 } else {
  for ( int i = 2 ; i <= n ; i += 2 )
   ret.push_back( i );
  for ( int i = 1 ; i <= n ; i += 2 )
   ret.push_back( i );
 return ret;
8.9 Aliens Optimization
long long Alien()
 long long c = kInf;
 for (int d = 60; d >= 0; --d) {
  // cost can be negative, depending on the problem.
  if (c - (1LL << d) < 0) continue;
  long long ck = c - (1LL \ll d)
```

pair<long long, int> r = check(ck);

```
if (r.second == k) return r.first - ck * k;
  if (r.second < k) c = ck;</pre>
pair<long long, int> r = check(c);
return r.first - c * k;
8.10 Hilbert Curve
long long hilbert(int n, int x, int y) {
long long res = 0;
for (int s = n / 2; s; s >>= 1) {
  int rx = (x & s) > 0, ry = (y & s) > 0;
  res += s * 111 * s * ((3 * rx) ^ ry);
 if (ry == 0) {
  if (rx == 1) x = s - 1 - x, y = s - 1 - y;
   swap(x, y);
  }
return res;
8.11 Binary Search On Fraction
struct Q {
11 p, q;
Q go(Q b, 11 d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 \le p,q \le N
Q frac_bs(11 N) {
Q lo{0, 1}, hi{1, 0};
if (pred(lo)) return lo;
 assert(pred(hi));
bool dir = 1, L = 1, H = 1;
for (; L || H; dir = !dir) {
  11 len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
   if (Q mid = hi.go(lo, len + step);
     mid.p > N || mid.q > N || dir ^ pred(mid))
   t++;
else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
return dir ? hi : lo;
```