Contents

1	Basi	C										1
•	1.1	vimrc										1
	1.2	Debug Macro										1
	1.3	Increase Stack										1
	1.4	Pragma Optimization										2
	1.5	IO Optimization									 •	2
2	Data	a Structure										2
_	2.1	Dark Magic										2
	2.2	Link-Cut Tree										2
	2.3	LiChao Segment Tree										3
	2.4	Treap										3
	2.5	Linear Basis										3
	2.6	Binary Search On Segment Tree	 •	•	 •		 ٠	٠	٠			3
3	Gra	oh										4
	3.1	2-SAT (SCC)										4
	3.2	BCC										4
	3.3	Centroid Decomposition										4
	3.4	Directed Minimum Spanning Tree .										5
	3.5	Dominator Tree										5
	3.6 3.7	Edge Coloring										5
	3.8	Manhattan Minimum Spanning Tree										6
	3.9	MaxClique										6
		MaxCliqueDyn										7
	3.11	Minimum Mean Cycle										7
		Minimum Steiner Tree										7
		Mo's Algorithm on Tree										8
	3.14	Virtual Tree	 ٠	•	 •	 ٠	 ٠	•	•	•		8
4	Mat	ching & Flow										8
•	4.1	Bipartite Matching										8
	4.2	Dijkstra Cost Flow										8
	4.3	Dinic										9
	4.4	Flow Models										9
	4.5	General Graph Matching										9
	4.6 4.7	Global Min-Cut										10 10
	4.7	GomoryHu Tree										10 10
	4.9	Minimum Cost Circulation										11
	4.10	Minimum Cost Maximum Flow										11
	4.11	Maximum Weight Graph Matching .										11
												47
_		L_										
5	Mat											13 13
5	Mat 5.1	Strling Number										13
5												
5		Strling Number									 	13 13
5	5.1 5.2 5.3	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey			 	 	 				 	13 13 13 13 13
5	5.1 5.2 5.3 5.4	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial	 		 	 	 				 	13 13 13 13 13
5	5.1 5.2 5.3 5.4 5.5	Strling Number	 		 	 	 				 	13 13 13 13 13 13
5	5.1 5.2 5.3 5.4 5.5 5.6	Strling Number			 	 	 				 	13 13 13 13 13 13 13
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7	Strling Number			 	 	 					13 13 13 13 13 13
5	5.1 5.2 5.3 5.4 5.5 5.6	Strling Number			 	 	 					13 13 13 13 13 13 13
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9	Strling Number			 	 	 					13 13 13 13 13 13 14 14
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10	Strling Number			 	 	 					13 13 13 13 13 13 14 14 14
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12	Strling Number 5.1.1 First Kind 5.1.2 Second Kind 5.1.2 Second Kind 0.2 Second Kind 0.3 Second Kind 0.4 Second Kind 0.5 Second			 	 	 					13 13 13 13 13 13 14 14 14 15 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Miller Rabin			 	 	 					13 13 13 13 13 13 14 14 14 15 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT			 	 						13 13 13 13 13 13 14 14 15 15 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number										13 13 13 13 13 13 14 14 15 15 15 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve)										13 13 13 13 13 13 14 14 15 15 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number										13 13 13 13 13 13 14 14 15 15 15 15 16
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho										13 13 13 13 13 13 14 14 15 15 15 16 16
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5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.15 5.14 5.15 5.16 5.17 5.18 5.19 5.19 5.20	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Polynomial Operations Quadratic residue										13 13 13 13 13 13 14 14 15 15 15 16 16 16 17
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.15 5.16 5.17 5.18 5.19 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.19 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.19 5.19 5.19 5.19 5.19 5.19 5.19	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction										13 13 13 13 13 13 14 14 15 15 15 16 16 17 17
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.15 5.16 5.17 5.18 5.19 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.19 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.19 5.19 5.19 5.19 5.19 5.19 5.19	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction										13 13 13 13 13 13 14 14 14 15 15 16 16 17 17
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.15 5.16 5.17 5.18 5.19 5.20 5.21 Geo	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction										13 13 13 13 13 13 14 14 14 15 15 16 16 17 17 17 18
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.15 5.16 5.17 5.18 5.19 5.20 5.21 Geo 6.1	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction metry Basic Geometry										13 13 13 13 13 14 14 15 15 15 16 16 17 17 18 18
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.15 5.16 5.17 5.18 5.19 5.20 6.2 6.3 6.4	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull										13 13 13 13 13 14 14 15 15 15 16 16 17 17 18 18 18 18
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	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.5 6.6 6.7	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair kD Closest Pair (3D ver.) Simulated Annealing										13 13 13 13 13 13 14 14 14 14 15 15 15 15 16 16 16 17 17 17 18 18 18 18 18 18 19
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.14 5.15 5.16 5.17 5.18 5.20 5.21 Geo 6.2 6.3 6.4 6.6 6.6 6.7 6.8	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection										13 13 13 13 13 13 13 13 14 14 14 15 15 15 15 16 16 16 17 17 18 18 18 18 19 19
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.15 5.16 5.17 5.18 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 2D Farthest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum										13 13 13 13 13 13 14 14 14 14 15 15 15 15 16 16 16 17 17 17 18 18 18 18 18 18 19
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.15 5.16 5.17 5.18 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection										13 13 13 13 13 13 13 13 14 14 14 15 15 15 15 16 16 17 17 18 18 18 18 18 19 19 19
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	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.14 5.15 5.16 5.17 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.10 6.11 6.12 6.13 6.14 6.15 6.14 6.15 6.16 6.17 6.17 6.17 6.17 6.17 6.17 6.17	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 2D Farthest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Circle Class Intersection of Polygon and Circle Point & Hulls Tangent Convex Hull Tangent										13 13 13 13 13 13 14 14 14 14 15 15 15 15 15 16 16 16 17 17 17 18 18 18 18 18 19 19 19 19 19 19 19 19 19 19 19 19 19

```
7 Stringology
  22
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24
1
Basic
1.1 vimrc
```

```
se is nu ru et tgc sc hls cin cino+=j1 sw=4 sts=4 bs=2
    mouse=a encoding=utf-8 ls=2
svn on
colorscheme desert
filetype indent on
inoremap {<CR> {<CR>}<ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -std=c++17 -
    DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
    Wconversion -fsanitize=address,undefined -g && echo
     success<CR>
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 &&
     echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
```

1.2 Debug Macro

```
#ifdef KISEKI
#define safe cerr<<__PRETTY_FUNCTION__\
    <<" line "<<__LINE__<<" safe\n"</pre>
#define debug(a...) gwerty(#a, a)
#define orange(a...) dvorak(#a, a)
using std::cerr;
template <typename ...T>
void qwerty(const char *s, T ...a) {
  cerr << "\e[1;32m(" << s << ") = (";</pre>
   int cnt = sizeof...(T);
   (..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
template <typename Iter>
void dvorak(const char *s, Iter L, Iter R) {
  cerr << "\e[1;32m[ " << s << " ] = [ ";</pre>
  for (int f = 0; L != R; ++L)

cerr << (f++ ? ", " : "") << *L;

cerr << " ]\e[0m\n";
#else
#define safe ((void)0)
#define debug(...) ((void)0)
#define orange(...) ((void)0)
#endif
```

1.3 Increase Stack

```
const int size = 256 << 20;</pre>
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));
```

1.4 Pragma Optimization

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,sse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8000)
;
```

1.5 IO Optimization

```
| static inline int gc() {
    constexpr int B = 1<<20;
    static char buf[B], *p, *q;
    if(p == q &&
        (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
        return EOF;
    return *p++;
    }
    template < typename T >
    static inline bool gn( T &x ) {
        int c = gc(); T sgn = 1; x = 0;
        while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
        if(c == '-') sgn = -1, c = gc();
        if(c == EOF) return false;
        while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
        return x *= sgn, true;
}
```

2 Data Structure

2.1 Dark Magic

2.2 Link-Cut Tree

```
template <typename Val> class LCT {
private:
struct node
  int pa, ch[2];
 bool rev;
 Val v, v_prod, v_rprod;
 node() \; : \; pa\{0\}, \; ch\{0, \; 0\}, \; rev\{false\}, \; v\{\}, \; v\_prod\{\}, \\
    v_rprod{} {};
};
vector<node> nodes;
 set<pair<int, int>> edges;
bool is_root(int u) const {
 const int p = nodes[u].pa;
 return nodes[p].ch[0] != u and nodes[p].ch[1] != u;
bool is_rch(int u) const {
 return (not is_root(u)) and nodes[nodes[u].pa].ch[1]
    == u;
void down(int u) {
 if (auto &cnode = nodes[u]; cnode.rev) {
   if (cnode.ch[0]) set_rev(cnode.ch[0]);
   if (cnode.ch[1]) set_rev(cnode.ch[1]);
   cnode.rev = false;
 void up(int u) {
 auto &cnode = nodes[u];
 cnode.v_prod =
  nodes[cnode.ch[0]].v_prod * cnode.v * nodes[cnode.ch
    [1]].v_prod;
  cnode.v_rprod =
   nodes[cnode.ch[1]].v_rprod * cnode.v * nodes[cnode.
    ch[0]].v_rprod;
```

```
void set_rev(int u) {
  swap(nodes[u].ch[0], nodes[u].ch[1]);
  swap(nodes[u].v_prod, nodes[u].v_rprod);
 nodes[u].rev ^= 1;
 void rotate(int u) {
 int f = nodes[u].pa, g = nodes[f].pa, l = is_rch(u);
if (nodes[u].ch[l ^ 1])
  nodes[nodes[u].ch[1 ^ 1]].pa = f;
  if (not is_root(f))
   nodes[g].ch[is_rch(f)] = u;
  nodes[f].ch[1] = nodes[u].ch[1 ^ 1];
  nodes[u].ch[1 ^ 1] = f
  nodes[u].pa = g, nodes[f].pa = u;
 up(f);
 void splay(int u) {
  vector<int> stk = {u};
  while (not is_root(stk.back()))
  stk.push_back(nodes[stk.back()].pa);
  for (; not stk.empty(); stk.pop_back())
   down(stk.back())
  for(int f=nodes[u].pa;!is_root(u);f=nodes[u].pa){
   if(!is_root(f))rotate(is_rch(u)==is_rch(f)?f:u);
   rotate(u);
  }
 up(u);
 void access(int u) {
  int last = 0;
  for (int last = 0; u; last = u, u = nodes[u].pa) {
   splay(u);
   nodes[u].ch[1] = last;
   up(u);
 int find_root(int u) {
  access(u); splay(u);
  int la = 0:
  for (; u; la = u, u = nodes[u].ch[0]) down(u);
  return la;
 void change_root(int u) {
 access(u); splay(u); set_rev(u);
 void link(int x, int y) {
 change_root(y); nodes[y].pa = x;
 void split(int x, int y) {
 change_root(x); access(y); splay(y);
 void cut(int x, int y) {
 split(x, y)
  nodes[y].ch[0] = nodes[x].pa = 0;
 up(y);
public:
LCT(int n = 0) : nodes(n + 1) {}
 int add(const Val &v = {}) {
 nodes.push_back(v)
  return int(nodes.size()) - 2;
 int add(Val &&v) {
 nodes.emplace_back(move(v));
  return int(nodes.size()) - 2;
 void set_val(int u, const Val &v) {
 splay(++u); nodes[u].v = v; up(u);
 Val query(int x, int y) {
 split(++x, ++y)
  return nodes[y].v_prod;
 bool connected(int u, int v) { return find_root(++u)
    == find_root(++v); }
 void add_edge(int u, int v) {
 if (++u > ++v) swap(u, v)
  edges.emplace(u, v); link(u, v);
 void del_edge(int u, int v) {
 auto k = minmax(++u, ++v);
```

```
if (auto it = edges.find(k); it != edges.end()) {
                                                                } // sz(L) == k
   edges.erase(it); cut(u, v);
                                                                int getRank(node *o) { // 1-base
 }
                                                                 int r = sz(o->lc) + 1;
};
                                                                 for (;o->pa != nullptr; o = o->pa)
                                                                  if (o->pa->rc == o) r += sz(o->pa->lc) + 1;
2.3 LiChao Segment Tree
                                                                 return r:
struct L {
int m, k, id;
L() : id(-1) {}
L(int a, int b, int c) : m(a), k(b), id(c) {}
                                                                #undef sz
                                                               2.5 Linear Basis
 int at(int x) { return m * x + k; }
                                                               template <int BITS> struct Basis {
class LiChao {
                                                                array<pair<uint64_t, int>, BITS> b;
                                                                Basis() { b.fill({0, -1}); }
private:
                                                                void add(uint64_t x, int p) {
  for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
 int n; vector<L> nodes;
static int lc(int x) { return 2 * x + 1;
static int rc(int x) { return 2 * x + 2;
                                                                  if (b[i].first == 0) {
 void insert(int 1, int r, int id, L ln) {
                                                                  b[i] = \{x, p\};
  int m = (1 + r) >> 1;
                                                                   return;
  if (nodes[id].id == -1) {
                                                                  } else if (b[i].second > p) {
   nodes[id] = ln;
                                                                   swap(b[i].first, x), swap(b[i].second, p);
   return:
                                                                  x ^= b[i].first;
  bool atLeft = nodes[id].at(1) < ln.at(1);</pre>
                                                                 }
  if (nodes[id].at(m) < ln.at(m)) {</pre>
  atLeft ^= 1
                                                                bool ok(uint64_t x, int p) {
   swap(nodes[id], ln);
                                                                 for (int i = 0; i < BITS; ++i)</pre>
                                                                  if (((x >> i) \& 1) \text{ and } b[i].second < p)
                                                                   x ^= b[i].first;
  if (r - l == 1) return;
  if (atLeft) insert(l, m, lc(id), ln);
                                                                 return x == 0;
  else insert(m, r, rc(id), ln);
                                                              };
 int query(int 1, int r, int id, int x) {
                                                               2.6 Binary Search On Segment Tree
 int ret = 0, m = (1 + r) >> 1;
if (nodes[id].id != -1)
                                                               // find_first = x -> minimal x s.t. check( [a, x) )
                                                               // find_last = x \rightarrow maximal x s.t. check([x, b))
  ret = nodes[id].at(x);
  if (r - 1 == 1) return ret;
                                                               template <typename C>
  if (x < m) return max(ret, query(1, m, lc(id), x));
return max(ret, query(m, r, rc(id), x));</pre>
                                                               int find_first(int 1, const C &check) {
                                                                if (1 >= n) return n + 1;
                                                                1 += sz;
                                                                for (int i = height; i > 0; i--)
public:
                                                                 propagate(1 >> i);
 LiChao(int n_{-}) : n(n_{-}), nodes(n * 4) {}
                                                                Monoid sum = identity;
 void insert(L ln) { insert(0, n, 0, ln); }
                                                                do {
 int query(int x) { return query(0, n, 0, x); }
                                                                 while ((1 & 1) == 0) 1 >>= 1;
                                                                 if (check(f(sum, data[1]))) {
                                                                  while (1 < sz) {
2.4 Treap
                                                                   propagate(1);
namespace Treap{
                                                                   1 <<= 1;
 #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                                   auto nxt = f(sum, data[1]);
 struct node{
                                                                   if (not check(nxt)) {
                                                                    sum = nxt;
  int size:
  uint32_t pri;
                                                                    1++:
 node *lc, *rc, *pa;
                                                                   }
  node():size(0),pri(rand()),lc(0),rc(0),pa(0){}
  void pull() {
                                                                  return 1 + 1 - sz;
  size = 1; pa = nullptr;
   if ( lc ) { size += lc->size; lc->pa = this; }
                                                                 sum = f(sum, data[1++]);
   if ( rc ) { size += rc->size; rc->pa = this; }
                                                                } while ((1 & -1) != 1);
  }
                                                                return n + 1;
 node* merge( node* L, node* R ) {
                                                               template <typename C>
  if ( not L or not R ) return L ? L : R;
                                                               int find_last(int r, const C &check) {
  if ( L->pri > R->pri ) {
                                                                if (r <= 0) return -1;</pre>
  L->rc = merge( L->rc, R ); L->pull();
                                                                r += sz;
   return L;
                                                                for (int i = height; i > 0; i--)
  } else {
                                                                 propagate((r - 1) >> i);
   R->lc = merge( L, R->lc ); R->pull();
                                                                Monoid sum = identity;
                                                                do {
   return R;
  }
 }
                                                                 while (r > 1 and (r & 1)) r >>= 1;
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                                 if (check(f(data[r], sum))) {
 if ( not rt ) L = R = nullptr;
                                                                  while (r < sz) {</pre>
  else if( sz( rt->lc ) + 1 <= k ) {
                                                                   propagate(r);
                                                                   r = (r << 1) + 1;
                                                                   auto nxt = f(data[r], sum);
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
   L->pull();
                                                                   if (not check(nxt)) {
  } else {
                                                                    sum = nxt;
                                                                    r--;
   split_by_size( rt->lc, k, L, R->lc );
   R->pull();
```

```
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                                                                 low[u] = min(low[u], dfn[v]);
   return r - sz;
                                                                } else {
                                                                ++ch, dfs(v, u);
  sum = f(data[r], sum);
 } while ((r & -r) != r);
                                                                 low[u] = min(low[u], low[v]);
return -1;
                                                                 if (low[v] > dfn[u])
                                                                 bridge[t] = true
                                                                if (low[v] >= dfn[u])
                                                                 ap[u] = true;
3
    Graph
3.1 2-SAT (SCC)
                                                              ap[u] &= (ch != 1 or u != f);
class TwoSat{
private:
                                                            public:
int n;
                                                             void init(int n_) {
vector<vector<int>> rG,G,sccs;
                                                              g.assign(n = n_, vector<pair<int, int>>());
vector<int> ord,idx;
                                                              low.assign(n, ecnt = 0);
vector<bool> vis,result;
                                                              dfn.assign(n, 0);
void dfs(int u){
                                                              ap.assign(n, false);
 vis[u]=true
  for(int v:G[u])
                                                             void add_edge(int u, int v) {
  if(!vis[v]) dfs(v);
                                                              g[u].emplace_back(v, ecnt);
 ord.push_back(u);
                                                              g[v].emplace_back(u, ecnt++);
void rdfs(int u){
                                                             void solve() {
 vis[u]=false;idx[u]=sccs.size()-1;
                                                              bridge.assign(ecnt, false);
  sccs.back().push_back(u);
                                                              for (int i = 0; i < n; ++i)
  if (not dfn[i]) dfs(i, i);</pre>
  for(int v:rG[u])
   if(vis[v])rdfs(v);
                                                             bool is_ap(int x) { return ap[x]; }
public:
                                                             bool is_bridge(int x) { return bridge[x]; }
void init(int n_){
 G.clear();G.resize(n=n_);
 rG.clear();rG.resize(n);
                                                            3.3 Centroid Decomposition
  sccs.clear();ord.clear();
                                                            struct Centroid {
 idx.resize(n);result.resize(n);
                                                             vector<vector<int64_t>> Dist;
                                                             vector<int> Parent, Depth;
 void add_edge(int u,int v){
                                                             vector<int64_t> Sub, Sub2;
 G[u].push_back(v);rG[v].push_back(u);
                                                             vector<int> Sz, Sz2;
                                                             Centroid(vector<vector<pair<int, int>>> g) {
 void orr(int x,int y){
                                                              int N = g.size()
 if ((x^y)==1)return
                                                              vector<bool> Vis(N);
  add_edge(x^1,y); add_edge(y^1,x);
                                                              vector<int> sz(N), mx(N);
                                                              vector<int> Path;
bool solve(){
                                                              Dist.resize(N)
 vis.clear();vis.resize(n);
                                                              Parent.resize(N);
 for(int i=0;i<n;++i)</pre>
                                                              Depth.resize(N)
  if(not vis[i])dfs(i);
                                                              auto DfsSz = [&](auto dfs, int x) -> void {
  reverse(ord.begin(),ord.end());
                                                               Vis[x] = true; sz[x] = 1; mx[x] = 0;
  for (int u:ord){
                                                                for (auto [u, w] : g[x]) {
   if(!vis[u])continue;
                                                                if (Vis[u]) continue;
  sccs.push_back(vector<int>());
                                                                dfs(dfs, u)
  rdfs(u);
                                                                sz[x] += sz[u]
                                                                mx[x] = max(mx[x], sz[u]);
 for(int i=0;i<n;i+=2)</pre>
  if(idx[i]==idx[i+1])
                                                               Path.push_back(x);
   return false;
  vector<bool> c(sccs.size());
                                                              auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
  for(size_t i=0;i<sccs.size();++i){</pre>
                                                                -> void {
   for(auto sij : sccs[i]){
                                                               Dist[x].push_back(D);Vis[x] = true;
    result[sij]=c[i];
                                                               for (auto [u, w] : g[x]) {
    c[idx[sij^1]]=!c[i];
                                                                if (Vis[u]) continue;
  }
                                                                dfs(dfs, u, D + w);
 }
 return true;
                                                              };
                                                              auto Dfs = [&]
bool get(int x){return result[x];}
                                                               (auto dfs, int x, int D = 0, int p = -1)->void {
int get_id(int x){return idx[x];}
                                                               Path.clear(); DfsSz(DfsSz, x);
int count(){return sccs.size();}
                                                               int M = Path.size();
} sat2;
                                                               int C = -1;
                                                               for (int u : Path) {
3.2 BCC
                                                                 if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
class BCC {
                                                                Vis[u] = false;
private:
                                                               DfsDist(DfsDist, C);
int n, ecnt;
                                                               for (int u : Path) Vis[u] = false;
vector<vector<pair<int, int>>> g;
vector<int> dfn, low;
                                                               Parent[C] = p; Vis[C] = true;
                                                               Depth[C] = D;
vector<bool> ap, bridge;
void dfs(int u, int f) {
                                                               for (auto [u, w] : g[C]) {
                                                                if (Vis[u]) continue;
 dfn[u] = low[u] = dfn[f] + 1;
  int ch = 0;
                                                                dfs(dfs, u, D + 1, C);
  for (auto [v, t] : g[u]) if (v != f) {
  if (dfn[v]) {
```

};

```
Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
 Sz.resize(N); Sz2.resize(N);
                                                                 assert(false);
                                                                }
                                                               } DMST;
void Mark(int v) {
 int x = v, z = -1;
                                                               3.5 Dominator Tree
  for (int i = Depth[v]; i >= 0; --i) {
   Sub[x] += Dist[v][i]; Sz[x]++;
                                                               namespace dominator {
                                                               vector<int> g[maxn], r[maxn], rdom[maxn];
  if (z != -1)
                                                               int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
    Sub2[z] += Dist[v][i];
    Sz2[z]++;
                                                               int dom[maxn], val[maxn], rp[maxn], tk;
                                                               void init(int n) {
                                                                // vertices are numbered from \theta to n-1
   z = x; x = Parent[x];
 }
                                                                fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
                                                                fill(fa, fa + n, -1); fill(val, val + n, -1);
                                                                fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
int64_t Query(int v) {
                                                                fill(dom, dom + n, -1); tk = 0; for (int i = 0; i < n; ++i) {
 int64_t res = 0;
  int x = v, z = -1
 for (int i = Depth[v]; i >= 0; --i) {
                                                                 g[i].clear(); r[i].clear(); rdom[i].clear();
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
  if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
  z = x; x = Parent[x];
                                                               void add_edge(int x, int y) { g[x].push_back(y); }
                                                               void dfs(int x) {
                                                                rev[dfn[x] = tk] = x;
 return res;
}
                                                                fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
                                                                for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
      Directed Minimum Spanning Tree
                                                                 r[dfn[u]].push_back(dfn[x]);
struct DirectedMST { // find maximum
struct Edge {
                                                               void merge(int x, int y) { fa[x] = y; }
  int u, v;
                                                               int find(int x, int c = 0) {
  int w;
                                                                if (fa[x] == x) return c ? -1 : x;
 Edge(int u, int v, int w) : u(u), v(v), w(w) {}
                                                                int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
vector<Edge> Edges;
void clear() { Edges.clear(); }
                                                                if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
 void addEdge(int a, int b, int w) { Edges.emplace_back
                                                                fa[x] = p;
                                                                return c ? p : val[x];
    (a, b, w); }
 int solve(int root, int n) {
                                                               vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
 vector<Edge> E = Edges;
  int ans = 0;
 while (true) {
                                                               // p[i] = -2 if i is unreachable from s
   // find best in edge
                                                                dfs(s):
                                                                for (int i = tk - 1; i >= 0; --i)
   vector<int> in(n, -inf), prv(n, -1);
                                                                 for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
                                                                 if (i) rdom[sdom[i]].push_back(i);
     in[e.v] = e.w;
                                                                 for (int &u : rdom[i]) {
     prv[e.v] = e.u;
                                                                  int p = find(u);
                                                                  if (sdom[p] == i) dom[u] = i;
   in[root] = 0;
                                                                  else dom[u] = p;
   prv[root] = -1;
                                                                 if (i) merge(i, rp[i]);
   for (int i = 0; i < n; i++)
   if (in[i] == -inf)
                                                                vector<int> p(n, -2); p[s] = -1;
     return -inf;
                                                                for (int i = 1; i < tk; ++i)
  if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];</pre>
   // find cycle
   int tot = 0;
                                                                for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
   vector < int > id(n, -1), vis(n, -1);
   for (int i = 0; i < n; i++) {
                                                                return p;
    ans += in[i];
                                                               }}
    for (int x = i; x != -1 && id[x] == -1; x = prv[x])
                                                               3.6 Edge Coloring
                                                               // max(d_u) + 1 edge coloring, time: O(NM) int C[kN][kN], G[kN][kN]; // 1-based, G: ans
     if (vis[x] == i) {
      for (int y = prv[x]; y != x; y = prv[y])
       id[y] = tot;
                                                               void clear(int N) {
                                                                for (int i = 0; i <= N; i++)
for (int j = 0; j <= N; j++)
      id[x] = tot++;
      break;
                                                                   C[i][j] = G[i][j] = 0;
     }
     vis[x] = i;
                                                               void solve(vector<pair<int, int>> &E, int N) {
                                                                int X[kN] = {}, a;
auto update = [&](int u) {
   if (!tot)
    return ans:
                                                                 for (X[u] = 1; C[u][X[u]]; X[u]++);
   for (int i = 0; i < n; i++)</pre>
    if (id[i] == -1)
                                                                auto color = [&](int u, int v, int c) {
                                                                 int p = G[u][v];
     id[i] = tot++;
                                                                 G[u][v] = G[v][u] = c;
   // shrink
   for (auto &e : E) {
                                                                 C[u][c] = v, C[v][c] = u;
                                                                 C[u][p] = C[v][p] = 0;
   if (id[e.u] != id[e.v])
     e.w -= in[e.v];
                                                                 if (p) X[u] = X[v] = p
    e.u = id[e.u], e.v = id[e.v];
                                                                 else update(u), update(v);
                                                                 return p;
   n = tot;
   root = id[root];
                                                                auto flip = [&](int u, int c1, int c2) {
```

```
int p = C[u][c1];
 swap(C[u][c1], C[u][c2]);
 if (p) G[u][p] = G[p][u] = c2;
 if (!C[u][c1]) X[u] = c1;
 if (!C[u][c2]) X[u] = c2;
 return p;
};
for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {</pre>
auto [u, v] = E[t];
int v0 = v, c = X[u], c0 = c, d;
 vector<pair<int, int>> L; int vst[kN] = {};
 while (!G[u][v0]) {
  L.emplace_back(v, d = X[v]);
  if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
  c = color(u, L[a].first, c);
else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
    color(u, L[a].first, L[a].second);
  else if (vst[d]) break
  else vst[d] = 1, v = C[u][d];
 if (!G[u][v0]) {
  for (; v; v = flip(v, c, d), swap(c, d));
if (C[u][c0]) { a = int(L.size()) - 1;
   while (--a >= 0 && L[a].second != c);
   for(;a>=0;a--)color(u,L[a].first,L[a].second);
  } else t--;
```

Lowbit Decomposition

```
class LBD {
 int timer, chains;
 vector<vector<int>> G;
 vector<int> tl, tr, chain, head, dep, pa;
// chains : number of chain
 // tl, tr[u] : subtree interval in the seq. of u
 // head[i] : head of the chain i
// chian[u] : chain id of the chain u is on
 void predfs(int u, int f) {
  dep[u] = dep[pa[u] = f] + 1;
  for (int v : G[u]) if (v != f) {
   predfs(v, u);
   if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
    chain[u] = chain[v];
  if (chain[u] == 0) chain[u] = ++chains;
 void dfschain(int u, int f) {
  tl[u] = timer++;
  if (head[chain[u]] == -1)
   head[chain[u]] = u;
  for (int v : G[u])
   if (v != f and chain[v] == chain[u])
    dfschain(v, u)
  for (int v : G[u])
   if (v != f and chain[v] != chain[u])
    dfschain(v, u);
  tr[u] = timer;
public:
 LBD(int n) : timer(0), chains(0), G(n), tl(n), tr(n),
 chain(n), head(n, -1), dep(n), pa(n) {}
void add_edge(int u, int v) {
 G[u].push_back(v); G[v].push_back(u);
 void decompose() { predfs(0, 0); dfschain(0, 0); }
 PII get_subtree(int u) { return {tl[u], tr[u]}; }
 vector<PII> get_path(int u, int v) {
  vector<PII> res;
  while (chain[u] != chain[v]) {
   if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
   swap(u, v);
int s = head[chain[u]];
   res.emplace_back(tl[s], tl[u] + 1);
   u = pa[s];
  if (dep[u] < dep[v]) swap(u, v);</pre>
  res.emplace_back(tl[v], tl[u] + 1);
  return res;
```

```
3.8 Manhattan Minimum Spanning Tree
```

```
typedef Point<int> P;
vector<array<int, 3>> manhattanMST(vector<P> ps) {
 vi id(sz(ps));
 iota(all(id), 0);
 vector<array<int, 3>> edges;
 rep(k, 0, 4) {
  sort(all(id), [&](int i, int j) {
   return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
  map<int, int> sweep;
  for (int i : id) {
   for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++)) {
    int j = it->second;
    P d = ps[i] - ps[j];
    if (d.y > d.x) break;
    edges.push_back({d.y + d.x, i, j});
   sweep[-ps[i].y] = i;
  for (P &p : ps)
   if (k \& 1) p.x = -p.x;
   else swap(p.x, p.y);
 return edges; // [{w, i, j}, ...]
}
```

3.9

};

```
MaxClique
// contain a self loop u to u, than u won't in clique
template < size_t MAXN >
class MaxClique{
private:
 using bits = bitset< MAXN >;
 bits popped, G[ MAXN ], ans;
 size_t deg[ MAXN ], deo[ MAXN ], n;
 void sort_by_degree() {
  popped.reset();
  for ( size_t i = 0 ; i < n ; ++ i )
    deg[ i ] = G[ i ].count();
  for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
    size_t mi = MAXN, id = 0;
    for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )</pre>
        mi = deg[id = j];
    popped[ deo[ i ] = id ] = 1;
for( size_t u = G[ i ]._Find_first() ;
     u < n ; u = G[ i ]._Find_next( u ) )</pre>
       -- deg[ u ];
  }
 void BK( bits R, bits P, bits X ) {
  if (R.count()+P.count() <= ans.count()) return;</pre>
  if ( not P.count() and not X.count() )
   if ( R.count() > ans.count() ) ans = R;
   return:
  /* greedily chosse max degree as pivot
  bits cur = P | X; size_t pivot = 0, sz = 0;
for ( size_t u = cur._Find_first() ;
   u < n ; u = cur._Find_next( u ) )</pre>
    if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
  cur = P & ( ~G[ pivot ] );
  */ // or simply choose first
  bits cur = P & (~G[ ( P | X )._Find_first() ]);
  for ( size_t u = cur._Find_first()
   u < n ; u = cur._Find_next( u ) ) {
   if ( R[ u ] ) continue;
   R[u] = 1;
   BK( R, P & G[ u ], X & G[ u ] );
   R[u] = P[u] = 0, X[u] = 1;
public:
 void init( size_t n_ ) {
  n = n_{-};
  for ( size_t i = 0 ; i < n ; ++ i )
   G[ i ].reset();
  ans.reset();
```

```
void add_edges( int u, bits S ) { G[ u ] = S; }
                                                                      d[i] = int((a[i] & mask).count());
void add_edge( int u, int v ) {
                                                                     sort(r.begin(), r.end(),
 G[u][v] = G[v][u] = 1;
                                                                      [&](int i, int j) { return d[i] > d[j]; });
                                                                     csort(r, c);
int solve() {
                                                                     dfs(r, c, 1, mask);
  sort_by_degree(); // or simply iota( deo... )
for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                                     return ans; // sol[0 ~ ans-1]
   deg[ i ] = G[ i ].count();
                                                                  } graph;
  bits pob, nob = 0; pob.set();
                                                                   3.11 Minimum Mean Cycle
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;
for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
                                                                  /* minimum mean cycle O(VE) */
   size_t v = deo[ i ];
                                                                   struct MMC{
  bits tmp; tmp[ v ] = 1;
BK( tmp, pob & G[ v ], nob & G[ v ] );
pob[ v ] = 0, nob[ v ] = 1;
                                                                   #define FZ(n) memset((n),0,sizeof(n))
                                                                   #define E 101010
                                                                   #define V 1021
                                                                  #define inf 1e9
  return static_cast< int >( ans.count() );
                                                                    struct Edge { int v,u; double c; };
                                                                    int n, m, prv[V][V], prve[V][V], vst[V];
                                                                    Edge e[E];
                                                                    vector<int> edgeID, cycle, rho;
3.10
      MaxCliqueDyn
                                                                    double d[V][V];
constexpr int kN = 150;
                                                                    void init( int _n ) { n = _n; m = 0; }
struct MaxClique { // Maximum Clique
bitset<kN> a[kN], cs[kN];
                                                                    // WARNING: TYPE matters
                                                                    void add_edge( int vi , int ui , double ci )
int ans, sol[kN], q, cur[kN], d[kN], n;
                                                                    { e[ m ++ ] = { vi , ui , ci }; }
void init(int _n) {
                                                                    void bellman_ford() {
                                                                     for(int i=0; i<n; i++) d[0][i]=0;</pre>
 n = _n, ans = q = 0;
  for (int i = 0; i < n; i++) a[i].reset();</pre>
                                                                     for(int i=0; i<n; i++) {</pre>
                                                                      fill(d[i+1], d[i+1]+n, inf);
for(int j=0; j<m; j++) {
  int v = e[j].v, u = e[j].u;
  if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
 void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
void csort(vector<int> &r, vector<int> &c) {
  int mx = 1, km = max(ans - q + 1, 1), t = 0,
    m = int(r.size())
                                                                        d[i+1][u] = d[i][v]+e[j].c;
                                                                        prv[i+1][u] = v;
  cs[1].reset(); cs[2].reset();
  for (int i = 0; i < m; i++) {
  int p = r[i], k = 1;
  while ((cs[k] & a[p]).count()) k++;</pre>
                                                                        prve[i+1][u] = j;
   if (k > mx) cs[++mx + 1].reset();
   cs[k][p] = 1;
                                                                    double solve(){
   if (k < km) r[t++] = p;
                                                                     // returns inf if no cycle, mmc otherwise
  c.resize(m);
                                                                     double mmc=inf;
  if (t) c[t - 1] = 0;
                                                                     int st = -1
  for (int k = km; k <= mx; k++) {</pre>
                                                                     bellman_ford();
   for (int p = int(cs[k]._Find_first());
                                                                     for(int i=0; i<n; i++) {</pre>
      p < kN; p = int(cs[k]._Find_next(p))) {
                                                                      double avg=-inf;
                                                                      for(int k=0; k<n; k++) {</pre>
    r[t] = p; c[t++] = k;
   }
                                                                       if(d[n][i]<inf-eps)</pre>
                                                                        avg=max(avg,(d[n][i]-d[k][i])/(n-k));
                                                                       else avg=max(avg,inf);
 void dfs(vector<int> &r, vector<int> &c, int 1,
  bitset<kN> mask) {
                                                                      if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
  while (!r.empty()) {
   int p = r.back(); r.pop_back();
                                                                     FZ(vst);edgeID.clear();cycle.clear();rho.clear();
   mask[p] = 0;
                                                                     for (int i=n; !vst[st]; st=prv[i--][st]) {
                                                                      vst[st]++:
   if (q + c.back() <= ans) return;</pre>
   cur[q++] = p;
                                                                      edgeID.PB(prve[i][st]);
   vector<int> nr, nc;
                                                                      rho.PB(st);
   bitset<kN> nmask = mask & a[p];
   for (int i : r)
                                                                     while (vst[st] != 2) {
    if (a[p][i]) nr.push_back(i);
                                                                      int v = rho.back(); rho.pop_back();
                                                                      cycle.PB(v);
   if (!nr.empty()) {
    if (1 < 4) {
                                                                      vst[v]++;
     for (int i : nr)
d[i] = int((a[i] & nmask).count());
                                                                     reverse(ALL(edgeID));
     sort(nr.begin(), nr.end(),
                                                                     edgeID.resize(SZ(cycle));
      [&](int x, int y)
                                                                     return mmc;
        return d[x] > d[y];
      });
                                                                  } mmc;
                                                                   3.12 Minimum Steiner Tree
   csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
                                                                  // Minimum Steiner Tree
    ans = q; copy(cur, cur + q, sol);
                                                                   // 0(V 3^T + V^2 2^T)
                                                                   struct SteinerTree {
   c.pop_back(); q--;
                                                                   #define V 33
  }
                                                                   #define T 8
                                                                   #define INF 1023456789
 int solve(bitset<kN> mask) { // vertex mask
                                                                    int n, dst[V][V], dp[1 << T][V], tdst[V];</pre>
 vector<int> r, c;
                                                                    void init(int _n) {
  for (int i = 0; i < n; i++)
                                                                     n = _n;
for (int i = 0; i < n; i++) {</pre>
   if (mask[i]) r.push_back(i);
  for (int i = 0; i < n; i++)</pre>
                                                                      for (int j = 0; j < n; j++)</pre>
```

```
dst[i][j] = INF * (i != j);
                                                            struct BipartiteMatching {
                                                             vector<int> X[N];
                                                             int fX[N], fY[N], n;
                                                             bitset<N> vis;
 void add_edge(int ui, int vi, int wi) {
 dst[ui][vi] = min(dst[ui][vi], wi);
                                                             bool dfs(int x)
                                                              for (auto i : X[x]) if (not vis[i]) {
  dst[vi][ui] = min(dst[vi][ui], wi);
                                                               vis[i] = true;
                                                               if (fY[i] == -1 || dfs(fY[i])) {
 void shortest_path() {
  for (int k = 0; k < n; k++)
                                                                fY[fX[x] = i] = x;
  for (int i = 0; i < n; i++)
                                                                return true:
   for (int j = 0; j < n; j++)
                                                               }
     dst[i][j] = min(dst[i][j], dst[i][k] + dst[k][j]);
                                                              return false;
 int solve(const vector<int> &ter) {
 int t = (int)ter.size();
                                                             void init(int n_, int m) {
  for (int i = 1; i < (1 << t); i++)
                                                              vis.reset();
   fill_n(dp[i], n, INF);
                                                              fill(X, X + (n = n_{-}), vector<int>());
  fill_n(dp[0], n, 0);
                                                              memset(fX, -1, sizeof(int) * n);
  for (int msk = 1; msk < (1 << t); msk++) {</pre>
                                                              memset(fY, -1, sizeof(int) * m);
  if (msk == (msk & (-msk))) {
    int who = __lg(msk);
                                                             void add_edge(int x, int y) { X[x].push_back(y); }
    for (int i = 0; i < n; i++)
                                                             int solve() { // return how many pair matched
     dp[msk][i] = dst[ter[who]][i];
                                                              int cnt = 0;
                                                              for (int i = 0; i < n; i++) {</pre>
    continue:
                                                               vis.reset();
   for (int i = 0; i < n; i++)
                                                               cnt += dfs(i);
    for (int submsk = (msk - 1) & msk; submsk; submsk =
     (submsk - 1) & msk)
                                                              return cnt;
     dp[msk][i] = min(dp[msk][i], dp[submsk][i] + dp[
                                                             }
    msk ^ submsk][i]);
                                                            };
   for (int i = 0; i < n; i++) {
                                                            4.2 Dijkstra Cost Flow
    tdst[i] = INF;
    for (int j = 0; j < n; j++)
                                                            // kN = #(vertices)
     tdst[i] = min(tdst[i], dp[msk][j] + dst[j][i]);
                                                            // MCMF.{Init, AddEdge, MincostMaxflow}
                                                            // MincostMaxflow(source, sink, flow_limit, &cost)
  copy_n(tdst, n, dp[msk]);
                                                            using Pii = pair<int, int>;
                                                            constexpr int kInf = 0x3f3f3f3f, kN = 500;
  int ans = INF;
  for (int i = 0; i < n; i++)
                                                            struct Edge {
  ans = min(ans, dp[(1 << t) - 1][i]);
                                                             int to, rev, cost, flow;
  return ans;
                                                            }:
                                                            struct MCMF { // 0-based
} solver;
                                                             int n{}, m{}, s{}, t{};
                                                             vector<Edge> graph[kN];
3.13
      Mo's Algorithm on Tree
                                                             // Larger range for relabeling
                                                             int64_t dis[kN] = {}, h[kN] = {};
dfs u:
                                                             int p[kN] = {};
push u
                                                             void Init(int nn) {
iterate subtree
                                                              n = nn;
push u
                                                              for (int i = 0; i < n; i++) graph[i].clear();</pre>
Let P = LCA(u, v) with St(u) <= St(v)
if (P == u) query[St(u), St(v)]
                                                             void AddEdge(int u, int v, int f, int c) {
else query[Ed(u), St(v)], query[St(P), St(P)]
                                                              graph[u].push_back({v,
                                                               static_cast<int>(graph[v].size()), c, f});
3.14
      Virtual Tree
                                                              graph[v].push_back(
vector<pair<int, int>> build(vector<int> vs, int r) {
                                                               {u, static_cast<int>(graph[u].size()) - 1,
                                                                -c, 0});
vector<pair<int, int>> res;
sort(vs.begin(), vs.end(), [](int i, int j) {
  return dfn[i] < dfn[j]; });</pre>
                                                             bool Dijkstra(int &max_flow, int64_t &cost) {
                                                              priority_queue<Pii, vector<Pii>, greater<>> pq;
 vector<int> s = {r};
for (int v : vs) if (v != r) {
                                                              fill_n(dis, n, kInf);
  if (int o = lca(v, s.back()); o != s.back()) {
                                                              dis[s] = 0
  while (s.size() >= 2)
                                                              pq.emplace(0, s);
    if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
                                                              while (!pq.empty()) {
                                                               auto u = pq.top();
    res.emplace_back(s[s.size() - 2], s.back());
    s.pop_back();
                                                               pq.pop();
                                                               int v = u.second;
                                                               if (dis[v] < u.first) continue;</pre>
   if (s.back() != o) {
                                                               for (auto &e : graph[v]) {
   res.emplace_back(s.back(), o);
    s.back() = o;
                                                                auto new_dis =
                                                                 dis[v] + e.cost + h[v] - h[e.to];
   }
                                                                if (e.flow > 0 && dis[e.to] > new_dis) {
  }
 s.push_back(v);
                                                                 dis[e.to] = new_dis;
                                                                 p[e.to] = e.rev;
for (size_t i = 1; i < s.size(); ++i)</pre>
                                                                 pq.emplace(dis[e.to], e.to);
 res.emplace_back(s[i - 1], s[i]);
                                                               }
 return res;
                                                              if (dis[t] == kInf) return false;
                                                              for (int i = 0; i < n; i++) h[i] += dis[i];</pre>
```

int d = max_flow;

for (int u = t; u != s;

4 Matching & Flow

Bipartite Matching

```
u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   d = min(d, graph[e.to][e.rev].flow);
 max_flow -= d;
  cost += int64_t(d) * h[t];
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   e.flow += d;
   graph[e.to][e.rev].flow -= d;
 return true:
int MincostMaxflow(
 int ss, int tt, int max_flow, int64_t &cost) {
  this->s = ss, this->t = tt;
 cost = 0;
 fill_n(h, n, 0);
 auto orig_max_flow = max_flow;
 while (Dijkstra(max_flow, cost) && max_flow) {}
  return orig_max_flow - max_flow;
};
```

Dinic

```
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
    int to, rev;
    Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
       int u = bfs.front(); bfs.pop();
       for (auto e: G[u]) {
         if (e.cap <= 0 or lv[e.to]!=-1) continue;
bfs.push(e.to); lv[e.to] = lv[u] + 1;</pre>
       }
    return lv[ed] != -1;
  Cap DFS(int u, Cap f){
  if (u == ed) return f;
    Cap ret = 0;
    for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
       auto &e = G[u][i];
       if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
       Cap nf = DFS(e.to, min(f, e.cap));
       ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
       if (f == 0) return ret;
    if (ret == 0) lv[u] = -1;
     return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
    G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
  st = st_, ed = ed_; Cap ret = 0;
    while (BFS()) {
       idx.assign(n, 0);
       Cap f = DFS(st, numeric_limits<Cap>::max());
       ret += f;
       if (f == 0) break;
     return ret;
};
```

4.4 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}.$
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t\to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f
 eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f'
 eq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipartite graph(X,Y)
 - 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect $x \to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y \to x$ with (cost,cap)=(-c,1)
 - 3. For each edge with c < 0, sum these cost as K, then increase d(y)by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v) > 0, connect $S \to v$ with (cost, cap) =(0, d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =(0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let ${\cal K}$ be the sum of all weights
 - 3. Connect source $s \to v$, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect u o v and v o u with capacity
 - 5. For $v \in {\it G}$, connect it with sink $v \to t$ with capacity K + 2T - $(\sum_{e \in E(v)} w(e)) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v
 - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with ca-
- 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4.5 General Graph Matching

```
namespace matching {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n)
 for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
 g[u].push_back(v);
 g[v].push_back(u);
int Find(int u) {
return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
static int tk = 0; tk++;
```

```
x = Find(x), y = Find(y);
                                                                 for (int i = 0; i < n - 1; ++i) {
 for (; ; swap(x, y)) {
  if (x != n) {
                                                                  int s, t; tie(s, t) = phase(n);
                                                                  del[t] = true; cut = min(cut, g[t]);
                                                                  for (int j = 0; j < n; ++j) {
  w[s][j] += w[t][j]; w[j][s] += w[j][t];</pre>
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
                                                                 }
                                                                 return cut;
void Blossom(int x, int y, int 1) {
  while (Find(x) != 1) {
                                                                4.7 GomoryHu Tree
 pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
                                                               int g[maxn];
                                                               vector<edge> GomoryHu(int n){
  if (fa[x] == x) fa[x] = 1;
                                                                 vector<edge> rt;
  if (fa[y] == y) fa[y] = 1;
                                                                 for(int i=1;i<=n;++i)g[i]=1;</pre>
                                                                 for(int i=2;i<=n;++i){</pre>
  x = pre[y];
                                                                  int t=g[i];
                                                                  flow.reset(); // clear flows on all edge
bool Bfs(int r, int n) {
                                                                  rt.push_back({i,t,flow(i,t)});
 for (int i = 0; i \le n; ++i) fa[i] = i, s[i] = -1;
                                                                  flow.walk(i); // bfs points that connected to i (use
                                                                    edges not fully flow)
 while (!q.empty()) q.pop();
 q.push(r);
                                                                  for(int j=i+1;j<=n;++j){</pre>
 s[r] = 0;
                                                                   if(g[j]==t && flow.connect(j))g[j]=i; // check if i
 while (!q.empty()) {
                                                                    can reach i
  int x = q.front(); q.pop();
  for (int u : g[x]) {
                                                                 }
   if (s[u] == -1) {
                                                                 return rt;
    pre[u] = x, s[u] = 1;
    if (match[u] == n) {
                                                                4.8 Kuhn Munkres
     for (int a = u, b = x, last; b != n; a = last, b =
                                                               class KM {
      last = match[b], match[b] = a, match[a] = b;
                                                               private:
     return true;
                                                                 static constexpr 11d INF = 1LL << 60;</pre>
                                                                 vector<lld> hl,hr,slk;
    q.push(match[u]);
                                                                 vector<int> fl,fr,pre,qu;
    s[match[u]] = 0;
                                                                 vector<vector<lld>> w;
   } else if (!s[u] && Find(u) != Find(x)) {
                                                                 vector<bool> v1,vr;
    int 1 = LCA(u, x, n);
                                                                 int n, ql, qr;
    Blossom(x, \dot{u}, \dot{l});
Blossom(\dot{u}, \dot{x}, \dot{l});
                                                                 bool check(int x) {
                                                                  if (v1[x] = true, f1[x] != -1)
                                                                   return vr[qu[qr++] = f1[x]] = true;
  }
                                                                  while (x != -1) swap(x, fr[fl[x] = pre[x]]);
                                                                  return false;
 return false;
                                                                 void bfs(int s) {
int Solve(int n) {
                                                                  fill(slk.begin(), slk.end(), INF);
 int res = 0;
                                                                  fill(vl.begin(), vl.end(), false);
 for (int x = 0; x < n; ++x) {
                                                                  fill(vr.begin(), vr.end(), false);
  if (match[x] == n) res += Bfs(x, n);
                                                                  ql = qr = 0;
                                                                  vr[qu[qr++] = s] = true;
 return res;
                                                                  while (true) {
}}
                                                                   11d d;
                                                                   while (ql < qr) {</pre>
4.6 Global Min-Cut
                                                                    for (int x = 0, y = qu[ql++]; x < n; ++x) {
const int maxn = 500 + 5
                                                                     if(!vl[x]&&slk[x]>=(d=hl[x]+hr[y]-w[x][y])){
int w[maxn][maxn], g[maxn];
                                                                      if (pre[x] = y, d) slk[x] = d;
bool v[maxn], del[maxn];
                                                                      else if (!check(x)) return;
void add_edge(int x, int y, int c) {
                                                                     }
w[x][y] += c; w[y][x] += c;
                                                                    }
                                                                   d = INF;
pair<int, int> phase(int n) {
                                                                   for (int x = 0; x < n; ++x)
memset(v, false, sizeof(v));
 memset(g, 0, sizeof(g));
                                                                    if (!vl[x] \&\& d > slk[x]) d = slk[x];
 int s = -1, t = -1;
                                                                   for (int x = 0; x < n; ++x) {
 while (true) {
                                                                    if (v1[x]) h1[x] += d;
                                                                    else slk[x] -= d;
  int c = -1:
  for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
                                                                    if (vr[x]) hr[x] -= d;
   if (c == -1 || g[i] > g[c]) c = i;
                                                                   for (int x = 0; x < n; ++x)
                                                                    if (!v1[x] && !slk[x] && !check(x)) return;
  if (c == -1) break;
                                                                  }
  v[s = t, t = c] = true;
  for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
                                                               public:
                                                                void init( int n_ ) {
   g[i] += w[c][i];
                                                                  qu.resize(n = n_);
  }
                                                                  fl.assign(n, -1); fr.assign(n, -1);
                                                                 hr.assign(n, 0); hl.resize(n);
 return make_pair(s, t);
                                                                  w.assign(n, vector<lld>(n));
                                                                  slk.resize(n); pre.resize(n);
int mincut(int n) {
                                                                  vl.resize(n); vr.resize(n);
 int cut = 1e9;
 memset(del, false, sizeof(del));
                                                                void set_edge( int u, int v, lld x ) {w[u][v] = x;}
```

```
11d solve() {
                                                              vector<Wei> dis;
  for (int i = 0; i < n; ++i)
                                                              PCW SPFA(){
  hl[i] = *max_element(w[i].begin(), w[i].end());
                                                               fill(inq.begin(),inq.end(),false);
                                                               fill(dis.begin(), dis.end(), INF_WEI);
  for (int i = 0; i < n; ++i) bfs(i);</pre>
 11d res = 0;
                                                               queue<int> qq; qq.push(ori);
 for (int i = 0; i < n; ++i) res += w[i][f1[i]];</pre>
                                                               dis[ori] = 0;
  return res;
                                                               while(not qq.empty()){
                                                                int u=qq.front();qq.pop();
} km;
                                                                inq[u] = false
                                                                for(int i=0;i<SZ(G[u]);++i){</pre>
4.9
      Minimum Cost Circulation
                                                                 Edge e=G[u][i];
struct Edge { int to, cap, rev, cost; };
                                                                 int v=e.to; Wei d=e.wei;
vector<Edge> g[kN];
                                                                 if(e.cap<=0||dis[v]<=dis[u]+d)</pre>
int dist[kN], pv[kN], ed[kN];
                                                                 dis[v] = dis[u] + d;
bool mark[kN];
                                                                 fa[v] = u, wh[v] = i;
if (inq[v]) continue;
int NegativeCycle(int n) {
memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
                                                                 qq.push(v);
int upd = -1;
                                                                 inq[v] = true;
for (int i = 0; i <= n; ++i) {
 for (int j = 0; j < n; ++j) {
                                                               if(dis[edd]==INF_WEI) return {-1, -1};
   int idx = 0;
   for (auto &e : g[j]) {
                                                               Cap mw=INF_CAP;
    if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
                                                               for(int i=edd;i!=ori;i=fa[i])
     dist[e.to] = dist[j] + e.cost;
                                                                mw=min(mw,G[fa[i]][wh[i]].cap);
                                                               for (int i=edd;i!=ori;i=fa[i]){
     pv[e.to] = j, ed[e.to] = idx;
                                                                auto &eg=G[fa[i]][wh[i]];
     if (i == n) {
      upd = j;
                                                                eg.cap -= mw;
      while(!mark[upd])mark[upd]=1,upd=pv[upd];
                                                                G[eg.to][eg.back].cap+=mw;
      return upd;
                                                               return {mw, dis[edd]};
    idx++;
                                                             public:
                                                              void init(int n){
                                                               G.clear();G.resize(n);
                                                               fa.resize(n);wh.resize(n);
return -1;
                                                               inq.resize(n); dis.resize(n);
int Solve(int n) {
                                                              void add_edge(int st, int ed, Cap c, Wei w){
int rt = -1, ans = 0;
                                                               G[st].emplace_back(ed,SZ(G[ed]),c,w);
while ((rt = NegativeCycle(n)) >= 0) {
                                                               G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
 memset(mark, false, sizeof(mark));
vector<pair<int, int>> cyc;
                                                              PCW solve(int a, int b){
 while (!mark[rt]) {
                                                               ori = a, edd = b;
                                                               Cap cc=0; Wei ww=0;
  cyc.emplace_back(pv[rt], ed[rt]);
  mark[rt] = true;
                                                               while(true)
   rt = pv[rt];
                                                                PCW ret=SPFA();
                                                                if(ret.first==-1) break;
  reverse(cyc.begin(), cyc.end());
                                                                cc+=ret.first;
                                                                ww+=ret.first * ret.second;
  int cap = kInf;
 for (auto &i : cyc)
  auto &e = g[i.first][i.second];
                                                               return {cc,ww};
                                                              }
   cap = min(cap, e.cap);
                                                             } mcmf;
  for (auto &i : cyc) {
                                                             4.11 Maximum Weight Graph Matching
  auto &e = g[i.first][i.second];
   e.cap -= cap;
                                                             struct WeightGraph {
  g[e.to][e.rev].cap += cap;
                                                              static const int inf = INT_MAX;
   ans += e.cost * cap;
                                                              static const int maxn = 514;
                                                              struct edge {
                                                               int u, v, w;
return ans;
                                                               edge(){}
                                                               edge(int u, int v, int w): u(u), v(v), w(w) {}
4.10 Minimum Cost Maximum Flow
                                                              int n, n_x;
class MiniCostMaxiFlow{
                                                              edge g[maxn * 2][maxn * 2];
using Cap = int; using Wei = int64_t;
                                                              int lab[maxn * 2];
using PCW = pair<Cap,Wei>;
                                                              int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa
static constexpr Cap INF_CAP = 1 << 30;</pre>
                                                                 [maxn * 2]:
static constexpr Wei INF_WEI = 1LL<<60;</pre>
                                                              int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
private:
                                                                 maxn * 2];
                                                              vector<int> flo[maxn * 2];
struct Edge{
                                                              queue<int> q;
 int to, back;
 Cap cap; Wei wei;
                                                              int e_delta(const edge &e) { return lab[e.u] + lab[e.v
                                                                 ] - g[e.u][e.v].w * 2; }
 Edge() {}
 Edge(int a,int b, Cap c, Wei d):
                                                              void update_slack(int u, int x) { if (!slack[x] ||
                                                                 e_{delta}(g[u][x]) < e_{delta}(g[slack[x]][x])) slack[x]
  to(a),back(b),cap(c),wei(d) {}
                                                                 ] = u; ]
int ori, edd;
                                                              void set_slack(int x) {
vector<vector<Edge>> G;
                                                               slack[x] = 0;
vector<int> fa, wh;
                                                               for (int u = 1; u <= n; ++u)
vector<bool> inq;
                                                                if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
```

```
update_slack(u, x);
void q_push(int x) {
 if (x \le n) q.push(x);
 else for (size_t i = 0; i < flo[x].size(); i++)</pre>
                                                                   xr):
   q_push(flo[x][i]);
void set_st(int x, int b) {
                                                                  pa[xs] = g[xns][xs].u;
 st[x] = b;
 if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)
    set_st(flo[x][i], b);
                                                                  q_push(xns);
int get_pr(int b, int xr) {
 int pr = find(flo[b].begin(), flo[b].end(), xr) - flo
   [b].begin();
                                                                  int xs = flo[b][i];
 if (pr % 2 == 1) {
  reverse(flo[b].begin() + 1, flo[b].end());
  return (int)flo[b].size() - pr;
                                                                 st[b] = 0;
 }
 return pr;
}
void set_match(int u, int v) {
                                                                 if (S[v] == -1)
                                                                  pa[v] = e.u, S[v] = 1;
 match[u] = g[u][v].v;
 if (u <= n) return;</pre>
                                                                  int nu = st[match[v]];
 edge e = g[u][v];
 int xr = flo_from[u][e.u], pr = get_pr(u, xr)
                                                                  S[nu] = 0, q_push(nu);
                                                                 } else if (S[v] == 0) {
 for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo</pre>
   [u][i ^ 1]);
 set_match(xr, v);
rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
   end());
                                                                 return false;
void augment(int u, int v) {
for (; ; ) {
  int xnv = st[match[u]];
                                                                bool matching() {
  set_match(u, v);
  if (!xnv) return
                                                                 q = queue<int>();
  set_match(xnv, st[pa[xnv]]);
u = st[pa[xnv]], v = xnv;
                                                                   q_push(x);
 }
                                                                 for (; ; ) {
  while (q.size()) {
int get_lca(int u, int v) {
 static int t = 0;
 for (++t; u || v; swap(u, v)) {
  if (u == 0) continue;
if (vis[u] == t) return u;
  vis[u] = t;
  u = st[match[u]];
  if (u) u = st[pa[u]];
 return 0;
                                                                    }
                                                                  int d = inf;
void add_blossom(int u, int lca, int v) {
 int b = n + 1;
 while (b <= n_x && st[b]) ++b;</pre>
 if (b > n_x) ++n_x
 lab[b] = 0, S[b] = 0;
 match[b] = match[lca];
 flo[b].clear();
 flo[b].push_back(lca);
 for (int x = u, y; x != lca; x = st[pa[y]])
                                                                   ]][x]) / 2);
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
 reverse(flo[b].begin() + 1, flo[b].end());
for (int x = v, y; x != lca; x = st[pa[y]])
                                                                   if (S[st[u]] == 0) {
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
                                                                    lab[u] -= d;
 set_st(b, b);
 for (int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].w
   = 0:
                                                                   if (st[b] == b) {
 for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
 for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  for (int x = 1; x <= n_x; ++x)
   if (g[b][x].w == 0 \mid \mid e_delta(g[xs][x]) < e_delta(g[xs][x])
                                                                  q = queue<int>();
   [b][x]))
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
  for (int x = 1; x <= n; ++x)
   if (flo_from[xs][x]) flo_from[b][x] = xs;
 set_slack(b);
```

```
void expand_blossom(int b) -
for (size_t i = 0; i < flo[b].size(); ++i)
  set_st(flo[b][i], flo[b][i]);</pre>
 int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
 for (int i = 0; i < pr; i += 2)
  int xs = flo[b][i], xns = flo[b][i + 1];
  S[xs] = 1, S[xns] = 0;
slack[xs] = 0, set_slack(xns);
 S[xr] = 1, pa[xr] = pa[b];
 for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
  S[xs] = -1, set_slack(xs);
bool on_found_edge(const edge &e) {
 int u = st[e.u], v = st[e.v];
  slack[v] = slack[nu] = 0;
  int lca = get_lca(u, v);
  if (!lca) return augment(u,v), augment(v,u), true;
  else add_blossom(u, lca, v);
memset(S + 1, -1, sizeof(int) * n_x);
memset(slack + 1, 0, sizeof(int) * n_x);
 for (int x = 1; x <= n_x; ++x)
 if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0,
 if (q.empty()) return false;
   int u = q.front(); q.pop();
   if (S[st[u]] == 1) continue;
   for (int v = 1; v <= n; ++v)
    if (g[u][v].w > 0 && st[u] != st[v]) {
     if (e_delta(g[u][v]) == 0) {
      if (on_found_edge(g[u][v])) return true;
     } else update_slack(u, st[v]);
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b && S[b] == 1) d = min(d, lab[b] / 2)
  for (int x = 1; x <= n_x; ++x)
   if (st[x] == x && slack[x]) {
    if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x])
    else if (S[x] == 0) d = min(d, e_delta(g[slack[x
  for (int u = 1; u <= n; ++u) {
    if (lab[u] <= d) return 0;</pre>
   } else if (S[st[u]] == 1) lab[u] += d;
  for (int b = n + 1; b \le n_x; ++b)
    if (S[st[b]] == 0) lab[b] += d * 2;
    else if (S[st[b]] == 1) lab[b] -= d * 2;
  for (int x = 1; x <= n_x; ++x)
   if (st[x] == x && slack[x] && st[slack[x]] != x &&
   e_delta(g[slack[x]][x]) == 0)
    if (on_found_edge(g[slack[x]][x])) return true;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b && S[b] == 1 && lab[b] == 0)
   expand_blossom(b);
```

```
return false;
 pair<long long, int> solve() {
 memset(match + 1, 0, sizeof(int) * n);
  n_x = n;
  int n_matches = 0;
  long long tot_weight = 0;
  for (int u = 0; u \le n; ++u) st[u] = u, flo[u].clear
    ():
  int w_max = 0;
  for (int u = 1; u <= n; ++u)
   for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);</pre>
    w_max = max(w_max, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
  while (matching()) ++n_matches;
  for (int u = 1; u <= n; ++u)
   if (match[u] && match[u] < u)</pre>
    tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
 void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
    g[vi][ui].w = wi; }
 void init(int _n) {
  n = _n;
  for (int u = 1; u <= n; ++u)
   for (int v = 1; v <= n; ++v)
    g[u][v] = edge(u, v, 0);
};
```

5 Math

5.1 Strling Number

5.1.1 First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1) \dots (x+n-1) = \sum_{k=0}^n S_1(n,k) x^k$$

$$g(x) = x(x+1) \dots (x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.1.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into knonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^{k} {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^{k} \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

5.2 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 \pmod{n}
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
if (y == 0) g=x,a=1,b=0;
else exgcd(y,x%y,g,b,a),b=(x/y)*a;
```

5.3 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
vector<T> d(output.size() + 1), me, he;
for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
  for (size_t j = 0; j < me.size(); ++j)
d[i] += output[i - j - 2] * me[j];</pre>
  if ((d[i] -= output[i - 1]) == 0) continue;
  if (me.empty()) {
   me.resize(f = i);
   continue;
```

```
vector<T> o(i - f - 1);
 T k = -d[i] / d[f]; o.push_back(-k);
 for (T x : he) o.push_back(x * k);
 if (o.size() < me.size()) o.resize(me.size());</pre>
 for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
 if (i-f+he.size() >= me.size()) he = me, f = i;
me = o:
return me;
```

5.4 Charateristic Polynomial

```
vector<vector<int>> Hessenberg(const vector<vector<int
    >> &A) {
 int N = A.size();
 vector<vector<int>> H = A;
 for (int i = 0; i < N - 2; ++i) {
  if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {
  if (H[j][i]) {</pre>
     for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
    ][k]);
     for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k]
    ][j]);
     break;
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
   int coef = 1LL * val * H[j][i] % kP;
for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL</pre>
     * H[i + 1][k] * (kP - coef)) % kP;
   for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
 return H;
vector<int> CharacteristicPoly(const vector<vector<int</pre>
    >> &A) {
 int N = A.size();
 auto H = Hessenberg(A);
 for (int i = 0; i < N; ++i) {
  for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];</pre>
 vector<vector<int>> P(N + 1, vector<int>(N + 1));
 P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
  P[i][0] = 0;
  for (int j = 1; j <= i; ++j) P[i][j] = P[i - 1][j -</pre>
  int val = 1;
  for (int j = i - 1; j >= 0; --j) {
   int coef = 1LL * val * H[i][i - 1] % kP:
   for (int k = 0; k \le j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 if (N & 1) {
  for (int i = 0; i \le N; ++i) P[N][i] = kP - P[N][i];
 return P[N];
```

Chinese Remainder

```
x = a1 \% m1
x = a2 \% m2
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
// 0 <= x < lcm(m1, m2)
5.6 De-Bruijn
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
```

```
National Taiwan University - ckiseki
  if (n \% p == 0)
   for (int i = 1; i <= p; ++i)
     res[sz++] = aux[i];
 } else {
  aux[t] = aux[t - p];
  db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {
   aux[t] = i;
    db(t + 1, t, n, k);
  }
int de_bruijn(int k, int n) {
 // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
 if (k == 1) {
  res[0] = 0;
  return 1;
 for (int i = 0; i < k * n; i++) aux[i] = 0;
 sz = 0:
 db(1, 1, n, k);
 return sz;
5.7 DiscreteLog
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
  // x^? \equiv y (mod M)
  Int t = 1, c = 0, g = 1;
  for (Int M_ = M; M_ > 0; M_ >>= 1)
    g = g * x % M;
  for (g = gcd(g, M); t % g != 0; ++c) {
    if (t == y) return c;
     t = t * x * M:
  if (y % g != 0) return -1;
  t /= g, y /= g, M /= g;
Int h = 0, gs = 1;
  for (; h * h < M; ++h) gs = gs * x % M;
unordered_map<Int, Int> bs;
  for (Int s = 0; s < h; bs[y] = ++s)
    y = y * x % M;
  for (Int s = 0; s < M; s += h) {
     t = t * gs % M;
     if (bs.count(t)) return c + s + h - bs[t];
  return -1;
5.8 Extended Euler
      a^b \equiv \begin{cases} a^b \mod \varphi(m) + \varphi(m) & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases}
5.9 ExtendedFloorSum
 g(a,b,c,n) = \sum_{i=1}^{n} i \lfloor \frac{ai+b}{c} \rfloor
                \left( \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \right)
               \begin{cases} +g(a \bmod c, b \bmod c, c, n), \\ 0, \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1)) \end{cases}
                                                            a \geq c \vee b \geq c
                -h(c, c-b-1, a, m-1)),
                                                            otherwise
```

$$\begin{split} g(a,b,c,n) &= \sum_{i=0} i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \end{cases} \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \\ h(a,b,c,n) &= \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{c} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \end{cases} \end{split}$$

nm(m+1) - 2g(c, c-b-1, a, m-1)

-2f(c, c-b-1, a, m-1) - f(a, b, c, n), otherwise

5.10 Fast Fourier Transform

```
const int mod = 1000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353;
const int M3 = 1004535809;
```

```
int superBigCRT(int64_t A, int64_t B, int64_t C) {
  static_assert (M1 <= M2 && M2 <= M3);
  constexpr int6\overset{\cdot}{4}_t r12 = modpow(M1, M2-2, M2);
  constexpr int64_t r13 = modpow(M1, M3-2, M3);
  constexpr int64_t r23 = modpow(M2, M3-2, M3);
  constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
  B = (B - A + M2) * r12 % M2;

C = (C - A + M3) * r13 % M3;
  C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
 for (int i = 0; i <= maxn; i++)</pre>
  omega[i] = cplx(cos(2 * pi * j / maxn),
     sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
 int z = __builtin_ctz(n) - 1;
 for (int i = 0; i < n; ++i) {
  int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^=(i >> j & 1)<<(z - j);
  if (x > i) swap(v[x], v[i]);
 for (int s = 2; s <= n; s <<= 1) {
  int z = s >> 1;
  for (int i = 0; i < n; i += s) {</pre>
   for (int k = 0; k < z; ++k) {
  cplx x = v[i + z + k] * omega[maxn / s * k];
    v[i + z + k] = v[i + k] - x;
    v[i+k] = v[i+k] + x;
void ifft(vector<cplx> &v, int n) {
 fft(v, n); reverse(v.begin() + 1, v.end());
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
VL convolution(const VI &a, const VI &b) {
 // Should be able to handle N <= 10^5, C <= 10^4
 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 vector<cplx> v(sz);
 for (int i = 0; i < sz; ++i) {
  double re = i < a.size() ? a[i] : 0;</pre>
  double im = i < b.size() ? b[i] : 0;</pre>
  v[i] = cplx(re, im);
 fft(v, sz);
 for (int i = 0; i <= sz / 2; ++i) {
  int j = (sz - i) & (sz - 1);
  cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
     * cplx(0, -0.25);
  if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
     ].conj()) * cplx(0, -0.25);
  v[i] = x;
 ifft(v, sz);
 VL c(sz);
 for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
 return c:
VI convolution_mod(const VI &a, const VI &b, int p) {
 int sz = 1;
 while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
 vector<cplx> fa(sz), fb(sz);
 for (int i = 0; i < (int)a.size(); ++i)</pre>
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 for (int i = 0; i < (int)b.size(); ++i)</pre>
  fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
 fft(fa, sz), fft(fb, sz);
 double r = 0.25 / sz;
 cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);

for (int i = 0; i <= (sz >> 1); ++i) {

  int j = (sz - i) & (sz - 1);
  cplx a1 = (fa[i] + fa[j].conj());
```

```
cplx a2 = (fa[i] - fa[j].conj()) * r2;
                                                                     if( x[ i ] >= MOD ) x[ i ] -= MOD;
  cplx b1 = (fb[i] + fb[j].conj()) * r3;
                                                                    if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
  cplx b2 = (fb[i] - fb[j].conj()) * r4;
                                                                   }
  if (i != j) {
  cplx c1 = (fa[j] + fa[i].conj());
                                                                 if( inv )
   cplx c2 = (fa[j] - fa[i].conj()) * r2;
                                                                  for( int i = 0 ; i < N ; i++ ) {
  x[ i ] *= inv( N, MOD );</pre>
   cplx d1 = (fb[j] + fb[i].conj()) * r3;
                                                                   x[ i ] %= MOD;
   cplx d2 = (fb[j] - fb[i].conj()) * r4;
   fa[i] = c1 * d1 + c2 * d2 * r5;
  fb[i] = c1 * d2 + c2 * d1;
                                                                      Miller Rabin
                                                                5.13
  fa[j] = a1 * b1 + a2 * b2 * r5;
 fb[j] = a1 * b2 + a2 * b1;
                                                                bool isprime(llu x) {
                                                                 static auto witn = [](llu a, llu u, llu n, int t) {
fft(fa, sz), fft(fb, sz);
                                                                  if (!a) return false;
vector<int> res(sz);
                                                                  while (t--) {
 for (int i = 0; i < sz; ++i) {
                                                                   11u a2 = mmul(a, a, n);
 long long a = round(fa[i].re), b = round(fb[i].re),
                                                                   if (a2 == 1 && a != 1 && a != n - 1) return true;
       c = round(fa[i].im);
                                                                   a = a2:
 res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
                                                                  return a != 1;
return res;
}}
                                                                 if (x < 2) return false;</pre>
                                                                 if (!(x & 1)) return x == 2;
5.11 FloorSum
                                                                 int t = __builtin_ctzll(x - 1);
// @param n `n < 2^32`
                                                                 llu odd = (x - 1) >> t;
// @param m `1 <= m < 2^32`
                                                                 for (llu m:
// @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
                                                                  {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
                                                                  if (witn(mpow(m % x, odd, x), odd, x, t))
11u ans = 0;
                                                                   return false;
while (true)
                                                                 return true;
 if (a >= m) {
   ans += n * (n - 1) / 2 * (a / m); a %= m;
                                                                5.14 NTT
 if (b >= m) {
                                                                template <int mod, int G, int maxn>
  ans += n * (b / m); b %= m;
                                                                struct NTT {
                                                                 static_assert (maxn == (maxn & -maxn));
 llu y_max = a * n + b;
                                                                 int roots[maxn];
 if (y_max < m) break;</pre>
                                                                 NTT () {
 // y_max < m * (n + 1)
                                                                  int r = modpow(G, (mod - 1) / maxn);
 // floor(y_max / m) <= n
                                                                  for (int i = maxn >> 1; i; i >>= 1) {
 n = (1lu)(y_max / m), b = (1lu)(y_max % m);
                                                                   roots[i] = 1;
 swap(m, a);
                                                                   for (int j = 1; j < i; j++)
roots[i + j] = modmul(roots[i + j - 1], r);</pre>
return ans;
                                                                   r = modmul(r, r);
11d floor_sum(lld n, lld m, lld a, lld b) {
11u \text{ ans} = 0:
                                                                 // n must be 2^k, and 0 \ll F[i] \ll mod
if (a < 0) {
                                                                 void operator()(int F[], int n, bool inv = false) {
 11u \ a2 = (a \% m + m) \% m;
                                                                  for (int i = 0, j = 0; i < n; i++) {
 ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
                                                                   if (i < j) swap(F[i], F[j]);</pre>
 a = a2:
                                                                   for (int k = n > 1; (j^k < k; k > = 1);
if (b < 0) {
                                                                  for (int s = 1; s < n; s *= 2) {
 11u b2 = (b \% m + m) \% m;
                                                                   for (int i = 0; i < n; i += s * 2) {
for (int j = 0; j < s; j++) {
 ans -= 1ULL * n * ((b2 - b) / m);
 b = b2;
                                                                     int a = F[i+j];
                                                                     int b = modmul(F[i+j+s], roots[s+j]);
return ans + floor_sum_unsigned(n, m, a, b);
                                                                     F[i+j] = modadd(a, b); // a + b
                                                                     F[i+j+s] = modsub(a, b); // a - b
5.12 FWT
                                                                    }
                                                                   }
/* xor convolution:
* x = (x0, x1) , y = (y0, y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                  if (inv) {
                                                                   int invn = modinv(n);
                                                                   for (int i = 0; i < n; i++)</pre>
* x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
* z = (1/2) * z''
                                                                    F[i] = modmul(F[i], invn);
                                                                   reverse(F + 1, F + n);
* or convolution:
* x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
 * and convolution:
                                                                NTT<2013265921, 31, 1048576> ntt;
* x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
const LL MOD = 1e9+7;
                                                                5.15 Partition Number
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
for( int d = 1 ; d < N ; d <<= 1 ) {</pre>
                                                                int b = sqrt(n);
                                                                ans[0] = tmp[0] = 1;
for (int i = 1; i <= b; i++) {
  int d2 = d << 1:
  for( int s = 0 ; s < N ; s += d2 )
   for( int i = s , j = s+d ; i < s+d ; i++, j++ ){</pre>
                                                                 for (int rep = 0; rep < 2; rep++)</pre>
                                                                 for (int j = i; j <= n - i * i; j++)
  modadd(tmp[j], tmp[j-i]);
for (int j = i * i; j <= n; j++)</pre>
    LL ta = x[ i ] , tb = x[ j ];
    x[ i ] = ta+tb;
    x[ j ] = ta-tb;
```

return 1u << (32 - __builtin_clz(n - 1));</pre>

```
modadd(ans[j], tmp[j - i * i]);
                                                               static NTT<mod,G,maxn> ntt; // coefficients in [0, P)
                                                               explicit Poly(int n = 1) : V(n) {}
5.16 Pi Count (Linear Sieve)
                                                               Poly(const V &v) : V(v) {}
                                                               Poly(const Poly &p, size_t n) : V(n) {
static constexpr int N = 1000000 + 5;
                                                                copy_n(p.data(), min(p.size(), n), data());
lld pi[N];
vector<int> primes;
                                                               Poly &irev() { return reverse(data(), data() + size())
bool sieved[N];
                                                                     *this; }
1ld cube_root(lld x){
                                                               Poly &isz(int sz) { return resize(sz), *this; }
1ld s=cbrt(x-static_cast<long double>(0.1));
                                                               Poly &iadd(const Poly &rhs) { // n() == rhs.n()
 while(s*s*s <= x) ++s;
                                                                fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
 return s-1;
                                                                return *this;
11d square_root(11d x){
                                                               Poly &imul(int k) {
 1ld s=sqrt(x-static_cast<long double>(0.1));
                                                                fi(0, size())(*this)[i] = modmul((*this)[i], k);
 while(s*s <= x) ++s;</pre>
                                                                return *this;
 return s-1;
                                                               Poly Mul(const Poly &rhs) const {
void init(){
                                                                const int sz = n2k(size() + rhs.size() - 1);
primes.reserve(N);
                                                                Poly X(*this, sz), Y(rhs, sz);
 primes.push_back(1);
                                                                ntt(X.data(), sz), ntt(Y.data(), sz);
 for(int i=2;i<N;i++) {</pre>
                                                                fi(0, sz) X[i] = modmul(X[i], Y[i]);
 if(!sieved[i]) primes.push_back(i);
                                                                ntt(X.data(), sz, true);
  pi[i] = !sieved[i] + pi[i-1];
                                                                return X.isz(size() + rhs.size() - 1);
  for(int p: primes) if(p > 1) {
  if(p * i >= N) break;
                                                               Poly Inv() const { // coef[0] != 0
   sieved[p * i] = true;
                                                                if (size() == 1) return V{modinv(*begin())};
   if(p % i == 0) break;
                                                                const int sz = n2k(size() * 2);
                                                                Poly X = Poly(*this, (size() + 1) / 2).Inv().isz(sz),
                                                                    Y(*this, sz);
                                                                \mathsf{ntt}(\hat{\mathsf{X}}.\mathsf{data}(),\ \mathsf{sz}),\ \mathsf{ntt}(\mathsf{Y}.\mathsf{data}(),\ \mathsf{sz});
11d phi(11d m, 11d n) {
                                                                fi(0, sz) X[i] = modmul(X[i], modsub(2, modmul(X[i],
 static constexpr int MM = 80000, NN = 500;
                                                                   Y[i])));
 static lld val[MM][NN];
                                                                ntt(X.data(), sz, true);
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
                                                                return X.isz(size());
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
                                                               Poly Sqrt() const { // coef[0] \in [1, mod)^2
 1ld ret = phi(m,n-1)-phi(m/primes[n],n-1);
                                                                if (size() == 1) return V{QuadraticResidue((*this)
 if(m < MM\&n < NN) val[m][n] = ret+1;
                                                                   [0], mod)};
 return ret;
                                                                Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
                                                                   size());
11d pi_count(11d);
                                                                return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
11d P2(11d m, 11d n) {
 11d sm = square_root(m), ret = 0;
 for(lld i = n+1;primes[i]<=sm;i++)</pre>
                                                               pair<Poly, Poly> DivMod(const Poly &rhs) const {
  ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
                                                                if (size() < rhs.size()) return {V{0}, *this};</pre>
 return ret;
                                                                const int sz = size() - rhs.size() + 1;
Poly X(rhs); X.irev().isz(sz);
11d pi_count(11d m) {
                                                                Poly Y(*this); Y.irev().isz(sz);
if(m < N) return pi[m];</pre>
                                                                Poly Q = Y.Mul(X.Inv()).isz(sz).irev();
11d n = pi_count(cube_root(m));
                                                                X = rhs.Mul(Q), Y = *this;
fi(0, size()) Y[i] = modsub(Y[i], X[i]);
return phi(m, n) + n - 1 - P2(m, n);
                                                                return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
5.17 Pollard Rho
                                                               Poly Dx() const {
// does not work when n is prime
                                                                Poly ret(size() - 1);
// return any non-trivial factor
                                                                fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
1lu pollard_rho(llu n) {
                                                                   1]);
 static auto f = [](llu x, llu k, llu m) {
                                                                 return ret.isz(max<int>(1, ret.size()));
    return add(k, mul(x, x, m), m); };
 if (!(n & 1)) return 2;
                                                               Poly Sx() const {
 mt19937 rnd(120821011);
                                                                Poly ret(size() + 1);
 while (true) {
                                                                fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
  llu y = 2, yy = y, x = rnd() % n, t = 1;
                                                                   this)[i]);
  for (llu sz = 2; t == 1; sz <<= 1, y = yy) {
  for (llu i = 0; t == 1 && i < sz; ++i) {
                                                                 return ret;
    yy = f(yy, x, n);
                                                               Poly Ln() const { // coef[0] == 1
    t = gcd(yy > y ? yy - y : y - yy, n);
                                                                return Dx().Mul(Inv()).Sx().isz(size());
   }
                                                               Poly Exp() const \{ // coef[0] == 0 \}
  if (t != 1 && t != n) return t;
                                                                if (size() == 1) return V{1};
                                                                Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size)
                                                                   ());
                                                                Poly Y = X.Ln(); Y[0] = mod - 1;
      Polynomial Operations
5.18
                                                                fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
                                                                return X.Mul(Y).isz(size());
using V = vector<int>
#define fi(1, r) for (int i = int(1); i < int(r); ++i)</pre>
template <int mod, int G, int maxn> struct Poly : V {
                                                               Poly Pow(const string &K) const {
                                                                int nz = 0;
 static uint32_t n2k(uint32_t n) {
  if (n <= 1) return 1;
                                                                while (nz < size() && !(*this)[nz]) ++nz;</pre>
```

int nk = 0, nk2 = 0;

```
for (char c : K) {
   nk = (nk * 10 + c - '0') % mod;
   nk2 = nk2 * 10 + c - '0';
   if (nk2 * nz >= size())
    return Poly(size());
   nk2 %= mod - 1;
  if (!nk && !nk2) return Poly(V{1}, size());
  Poly X = V(data() + nz, data() + size() - nz * (nk2 -
     1));
  int x0 = X[0];
  return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
    modpow(x0, nk2)).irev().isz(size()).irev();
 Poly InvMod(int L) { // (to evaluate linear recursion)
 Poly R{1, 0}; // *this * R mod x^L = 1 (*this[0] ==
  for (int level = 0; (1 << level) < L; ++level) {</pre>
   Poly 0 = R.Mul(Poly(data(), min<int>(2 << level,</pre>
    size())));
   Poly Q(2 \ll level); Q[0] = 1;
   for (int j = (1 << level); j < (2 << level); ++j)</pre>
    Q[j] = modsub(mod, O[j]);
   R = R.Mul(Q).isz(4 << level);
  return R.isz(L);
 }
 static int LinearRecursion(const V &a, const V &c,
    int64_t n) { // a_n = \sum c_j a_(n-j)}
  const int k = (int)a.size();
  assert((int)c.size() == k + 1);
 Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};

fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
  C[k] = 1;
  while (n) {
   if (n % 2) W = W.Mul(M).DivMod(C).second;
   n /= 2, M = M.Mul(M).DivMod(C).second;
  int ret = 0:
  fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
  return ret;
};
#undef fi
using Poly_t = Poly<998244353, 3, 1 << 20>;
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
5.19 Quadratic residue
struct S {
 int MOD, w;
 int64_t x, y;
 S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
  : MOD(m), w(w_{-}), x(x_{-}), y(y_{-}) {}
 S operator*(const S &rhs) const {
  int w_{-} = w;
  if (w_ == -1) w_ = rhs.w;
  assert(w_! = -1 \text{ and } w_ == rhs.w);
  return { MOD, w_,
 (x * rhs.x + y * rhs.y % MOD * w) % MOD,
   (x * rhs.y + y * rhs.x) % MOD };
int get_root(int n, int P) {
  if (P == 2 \text{ or } n == 0) \text{ return } n;
  if (qpow(n, (P - 1) / 2, P) != 1) return -1;
  auto check = [&](int x) {
  return qpow(x, (P - 1) / 2, P); };
if (check(n) == P-1) return -1;
  int64_t a; int w; mt19937 rnd(7122);
  do { a = rnd() % P;
    w = ((a * a - n) % P + P) % P;
  } while (check(w) != P - 1);
  return qpow(S(P, w, a, 1), (P + 1) / 2).x;
5.20 Simplex
namespace simplex {
// maximize c^Tx under Ax <= B</pre>
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
```

```
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
 double inv = 1.0 / d[r][s];
 for (int i = 0; i < m + 2; ++i)
  for (int j = 0; j < n + 2; ++j)
if (i != r && j != s)
d[i][j] -= d[r][j] * d[i][s] * inv;</pre>
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;</pre>
 for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;
 d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
 int x = m + z;
 while (true) {
  int s = -1;
  for (int i = 0; i <= n; ++i) {
   if (!z && q[i] == -1) continue;
   if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
  if (d[x][s] > -eps) return true;
  for (int i = 0; i < m; ++i) {
   if (d[i][s] < eps) continue;</pre>
    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
  pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
 m = b.size(), n = c.size();
 d = VVD(m + 2, VD(n + 2))
 for (int i = 0; i < m; ++i)
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)</pre>
  p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
 for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
 q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
 for (int i = 1; i < m; ++i)</pre>
  if (d[i][n + 1] < d[r][n + 1]) r = i;
 if (d[r][n + 1] < -eps) {</pre>
  pivot(r, n)
  if (!phase(1) || d[m + 1][n + 1] < -eps)</pre>
   return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
   int`s = min_element(d[i].begin(), d[i].end() - 1)
        - d[i].begin();
   pivot(i, s);
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)</pre>
  if (p[i] < n) \times [p[i]] = d[i][n + 1];
 return x;
5.21
       Simplex Construction
Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that for all 1 \leq j \leq m,
\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j and x_i \geq 0 for all 1 \leq i \leq n.
```

- 1. In case of minimization, let $c'_i = -c_i$
- 2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- $3. \sum_{1 \le i \le n} A_{ji} x_i = b_j$
 - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$
 - $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6 Geometry

6.1 Basic Geometry

```
#define IM imag
#define RE real
using 1ld = int64_t;
using llf = long double;
using PT = std::complex<lld>;
using PTF = std::complex<llf>
auto toPTF(PT p) { return PTF{RE(p), IM(p)}; }
int sgn(lld x) { return (x > 0) - (x < 0); }
lld dot(PT a, PT b) { return RE(conj(a) * b); }
lld cross(PT a, PT b) { return IM(conj(a) * b); }
int ori(PT a, PT b, PT c) {
return sgn(cross(b - a, c - a));
bool operator<(const PT &a, const PT &b) {</pre>
 return RE(a) != RE(b) ? RE(a) < RE(b) : IM(a) < IM(b);
int quad(PT p) {
 return (IM(p) == 0) // use sgn for PTF
  ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
int argCmp(PT a, PT b) {
  // -1 / 0 / 1 <-> < / == / > (atan2)
 int qa = quad(a), qb = quad(b);
 if (qa != qb) return sgn(qa - qb);
 return sgn(cross(b, a));
template <typename V> llf area(const V & pt) {
 11d ret = 0:
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
 return ret / 2.0;
PT rot90(PT p) { return PT{-IM(p), RE(p)}; }
PTF project(PTF p, PTF q) { // p onto q
 return dot(p, q) * q / dot(q, q);
11f FMOD(11f x) {
 if (x < -PI) x += PI * 2;
 if (x > PI) x -= PI * 2;
 return x;
```

6.2 Segment & Line Intersection

```
struct Segment {
 PT st, dir; // represent st + t*dir for 0<=t<=1
 Segment(PT's, PT'e) : st(s), dir(e - s) {}
 static bool valid(lld p, lld q) {
  // is there t s.t. 0 <= t <= 1 && qt == p ?
  if (q < 0) q = -q, p = -p;
  return 0 <= p && p <= q;
bool isInter(Segment A, PT P) {
 if (A.dir == PT(0)) return P == A.st;
 return cross(P - A.st, A.dir) == 0 &&
  Segment::valid(dot(P - A.st, A.dir), norm(A.dir));
template <typename U, typename V>
bool isInter(U A, V B) {
  if (cross(A.dir, B.dir) == 0)
  return // handle parallel yourself
   isInter(A, B.st) || isInter(A, B.st+B.dir) ||
 isInter(B, A.st) || isInter(B, A.st+A.dir);
PT D = B.st - A.st;
 11d C = cross(A.dir, B.dir);
 return U::valid(cross(D, A.dir), C) &&
   V::valid(cross(D, B.dir), C);
struct Line {
PT st, ed, dir;
 Line (PT s, PT e)
  : st(s), ed(e), dir(e - s) {}
PTF intersect(const Line &A, const Line &B) {
11f t = cross(B.st - A.st, B.dir) /
  llf(cross(A.dir, B.dir));
 return toPTF(A.st) + PTF(t) * toPTF(A.dir);
```

```
6.3 2D Convex Hull
// returns a convex hull in counterclockwise order
// for a non-strict one, change cross >= to >
vector<PT> convex_hull(vector<PT> p) {
 sort(all(p));
 if (p[0] == p.back()) return {p[0]};
 int n = p.size(), t = 0;
 vector<PT> h(n + 1);
           _ = 2, s = 0; _--; s = --t, reverse(all(p)))
 for (int _
  for (PT i : p) {
   while (t > s + 1 \&\& cross(i, h[t-1], h[t-2]) >= 0)
    t--:
   h[t++] = i;
  }
 return h.resize(t), h;
6.4 3D Convex Hull
// return the faces with pt indexes
int flag[MXN][MXN];
struct Point{
 ld x, y, z;
 Point operator * (const ld &b) const {
  return (Point) {x*b,y*b,z*b};}
 Point operator * (const Point &b) const {
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
Point ver(Point a, Point b, Point c) {
return (b - a) * (c - a);}
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0
 REP(i,n) REP(j,n) flag[i][j] = 0;
 vector<Face> now;
 now.emplace_back(0,1,2);
 now.emplace_back(2,1,0)
 for (int i=3; i<n; i++){
  ftop++; vector<Face> next;
REP(j, SZ(now)) {
   Face& f=now[j]; int ff = 0;
   {\tt ld d=(pt[i]-pt[f.a]).dot(}\\
     ver(pt[f.a], pt[f.b], pt[f.c]));
   if (d <= 0) next.push_back(f);</pre>
   if (d > 0) ff=ftop;
   else if (d < 0) ff=-ftop;</pre>
   flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
  REP(j, SZ(now)) {
  Face& f=now[j];
   if (flag[f.a][f.b] > 0 &&
     flag[f.a][f.b] != flag[f.b][f.a])
    next.emplace_back(f.a,f.b,i);
   if (flag[f.b][f.c] > 0 &&
     flag[f.b][f.c] != flag[f.c][f.b])
    next.emplace_back(f.b,f.c,i);
   if (flag[f.c][f.a] > 0 &&
     flag[f.c][f.a] != flag[f.a][f.c])
    next.emplace_back(f.c,f.a,i);
  now=next;
 return now;
6.5 2D Farthest Pair
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++)</pre>
 while(abs(cross(stk[i+1]-stk[i],
```

```
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {
  while(abs(cross(stk[i+1]-stk[i],
    stk[(pos+1)%n]-stk[i])) >
    abs(cross(stk[i+1]-stk[i],
    stk[pos]-stk[i]))) pos = (pos+1)%n;
ans = max({ans, dis(stk[i], stk[pos]),
    dis(stk[i+1], stk[pos])});
}
```

6.6 kD Closest Pair (3D ver.)

```
llf solve(vector<P> v) {
  shuffle(v.begin(), v.end(), mt19937());
  unordered_map<lld, unordered_map<lld,</pre>
```

```
unordered_map<lld, int>>> m;
                                                                  POP(lines[i], lines[i]);
                                                                  pt.push_back(intersect(que.back(), lines[i]));
 llf d = dis(v[0], v[1]);
 auto Idx = [&d] (11f x) -> 11d {
                                                                  que.push_back(lines[i]);
  return round(x * 2 / d) + 0.1;
 auto rebuild_m = [&m, &v, &Idx](int k) {
                                                                POP(que.front(), que.back())
  m.clear();
                                                                if (que.size() <= 1 ||</pre>
                                                                  argCmp(que.front().dir, que.back().dir) == 0)
  for (int i = 0; i < k; ++i)
   m[Idx(v[i].x)][Idx(v[i].y)]
                                                                  return 0:
    [Idx(v[i].z)] = i;
                                                                pt.push_back(intersect(que.front(), que.back()));
 }; rebuild_m(2);
                                                                return area(pt);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
     kz = Idx(v[i].z); bool found = false;
                                                             6.9 Minkowski Sum
  for (int dx = -2; dx <= 2; ++dx) {
                                                             vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
   const 11d nx = dx + kx;
                                                               hull(A), hull(B);
   if (m.find(nx) == m.end()) continue;
                                                               vector<pll> C(1, A[0] + B[0]), s1, s2;
   auto& mm = m[nx];
                                                               for(int i = 0; i < SZ(A); ++i)
s1.pb(A[(i + 1) % SZ(A)] - A[i]);</pre>
   for (int dy = -2; dy <= 2; ++dy) {
    const 11d ny = dy + ky;
                                                               for(int i = 0; i < SZ(B); i++)
    if (mm.find(ny) == mm.end()) continue;
                                                                s2.pb(B[(i + 1) % SZ(B)] - B[i]);
    auto& mmm = mm[ny];
                                                               for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
    for (int dz = -2; dz <= 2; ++dz) {
                                                                if (p2 >= SZ(B)
     const 11d nz = dz + kz;
                                                                  || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
     if (mmm.find(nz) == mmm.end()) continue;
                                                                 C.pb(C.back() + s1[p1++]);
     const int p = mmm[nz];
     if (dis(v[p], v[i]) < d) {
  d = dis(v[p], v[i]);</pre>
                                                                 C.pb(C.back() + s2[p2++]);
                                                               return hull(C), C;
      found = true;
     }
                                                             6.10 Circle Class
                                                             struct Circle { PTF o; llf r; };
  if (found) rebuild_m(i + 1);
                                                             vector<llf> intersectAngle(Circle A, Circle B) {
  else m[kx][ky][kz] = i;
                                                              PTF dir = B.o - A.o; llf d2 = norm(dir);
                                                               if (norm(A.r - B.r) >= d2) // norm(x) := |x|^2
 return d;
                                                                if (A.r < B.r) return {-PI, PI}; // A in B</pre>
                                                                else return {}; // B in A
                                                               if (norm(A.r + B.r) <= d2) return {};</pre>
      Simulated Annealing
                                                               11f dis = abs(dir), theta = arg(dir);
11f anneal() {
                                                               11f phi = acos((A.r * A.r + d2 - B.r * B.r) /
 mt19937 rnd_engine( seed );
                                                                 (2 * A.r * dis));
 uniform_real_distribution< llf > rnd( 0, 1 );
                                                               11f L = FMOD(theta - phi), R = FMOD(theta + phi);
 const llf dT = 0.001;
                                                               return { L, R };
 // Argument p
 11f S_cur = calc( p ), S_best = S_cur;
for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
                                                             vector<PTF> intersectPoint(Circle a, Circle b) {
 // Modify p to p_prime
                                                               llf d = abs(a.o - b.o);
  const llf S_prime = calc( p_prime );
                                                               if (d >= b.r+a.r || d <= abs(b.r-a.r)) return {};</pre>
  const 11f delta_c = S_prime - S_cur
                                                               11f dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
  llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
                                                               PTF dir = (a.o - b.o) / d;
  if ( rnd( rnd_engine ) <= prob )</pre>
                                                               PTF u = dir*d1 + b.o;
   S_cur = S_prime, p = p_prime;
                                                               PTF v = rot90(dir) * sqrt(max<llf>(0, b.r*b.r-d1*d1));
  if ( S_prime < S_best ) // find min</pre>
                                                               return \{u + v, u - v\};
   S_best = S_prime, p_best = p_prime;
return S_best;
                                                              6.11 Intersection of line and Circle
                                                             vector<PTF> line_interCircle(const PTF &p1,
6.8 Half Plane Intersection
                                                                const PTF &p2, const PTF &c, const double r)
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
bool operator<(const Line &lhs, const Line &rhs) {</pre>
                                                               PTF ft = p1 + project(c-p1, p2-p1), vec = p2-p1;
                                                               11f dis = abs(c - ft);
                                                               if (abs(dis - r) < eps) return {ft};</pre>
  if (int cmp = argCmp(lhs.dir, rhs.dir))
                                                               if (dis > r) return {};
    return cmp == -1;
                                                               vec = vec * sqrt(r * r - dis * dis) / abs(vec);
  return ori(lhs.st, lhs.ed, rhs.st) < 0;</pre>
                                                               return {ft + vec, ft - vec};
// intersect function is in "Segment Intersect"
                                                             6.12 Intersection of Polygon and Circle
11f HPI(vector<Line> &lines) {
                                                             // Divides into multiple triangle, and sum up
  sort(lines.begin(), lines.end());
                                                             // test by HDU2892
  deque<Line> que;
  deque<PTF> pt;
                                                             11f _area(PTF pa, PTF pb, llf r)
                                                               if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
  que.push_back(lines[0]);
  for (int i = 1; i < (int)lines.size(); i++) {</pre>
                                                               if (abs(pb) < eps) return 0;</pre>
                                                               11f S, h, theta;
    if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
                                                               llf a = abs(pb), b = abs(pa), c = abs(pb - pa);
#define POP(L, R) \
                                                               11f cosB = dot(pb, pb - pa) / a / c, B = acos(cosB);
                                                               llf cosC = dot(pa, pb) / a / b, C = acos(cosC);
    while (pt.size() > 0 \
      && ori(L.st, L.ed, pt.back()) < 0) \
                                                               if (a > r) ·
      pt.pop_back(), que.pop_back(); \
                                                               S = (C / 2) * r * r;
                                                                h = a * b * sin(C) / c
    while (pt.size() > 0 \
      && ori(R.st, R.ed, pt.front()) < 0) \
                                                                if (h < r && B < PI / 2)
      pt.pop_front(), que.pop_front();
                                                                 S = (acos(h / r) * r * r - h * sqrt(r*r - h*h));
```

```
} else if (b > r) {
  theta = PI - B - asin(sin(B) / r * a);
 S = 0.5 * a*r * sin(theta) + (C - theta) / 2 * r*r;
 } else
 S = 0.5 * sin(C) * a * b;
return S;
11f area_poly_circle(const vector<PTF> &poly,
 const PTF &0, const llf r) {
                                                                else
 11f S = 0;
for (int i = 0, N = poly.size(); i < N; ++i)</pre>
 S += _area(poly[i] - 0, poly[(i + 1) % N] - 0, r) *
    ori(0, poly[i], poly[(i + 1) % N]);
 return fabs(S);
6.13 Point & Hulls Tangent
#define above(P, Vi, Vj) (ori(P, Vi, Vj) > 0) // true
                                                                else
    if Vi is above Vj
#define below(P, Vi, Vj) (ori(P, Vi, Vj) < 0) // true</pre>
    if Vi is below Vj
// Rtangent_PointPolyC(): binary search for convex
                                                             }
    polygon right tangent
    Input: P = a 2D point (exterior to the polygon)
                                                            6.14
        n = number of polygon vertices
        V = array of vertices for a 2D convex polygon
11
    with V[n] = V[0]
// Return: index "i" of rightmost tangent point V[i]
int Rtangent_PointPolyC(PT P, int n, PT *V) {
                                                           //
int a, b, c;
int upA, dnC;
                                                            11
                                                           //
if (below(P, V[1], V[0]) && !above(P, V[n - 1], V[0]))
                                                            11
 return 0;
for (a = 0, b = n;;) {
 c = (a + b) / 2
  dnC = below(P, V[c + 1], V[c]);
  if (dnC && !above(P, V[c - 1], V[c]))
  return c;
 upA = above(P, V[a + 1], V[a]);
  if (upA) {
  if (dnC) {
   b = c;
   } else {
    if (above(P, V[a], V[c]))
     b = c;
    else
     a = c:
   }
  } else {
   if (!dnC) {
   a = c;
   } else {
    if (below(P, V[a], V[c]))
                                                               --ix2:
    b = c;
   else
                                                              }
     a = c;
  }
                                                             return;
}
                                                            6.15
// Ltangent_PointPolyC(): binary search for convex
    polygon left tangent
    Input: P = a 2D point (exterior to the polygon)
        n = number of polygon vertices
        V = array of vertices for a 2D convex polygon
//
    with V[n]=V[0]
   Return: index
                   "i" of leftmost tangent point V[i]
int Ltangent_PointPolyC(PT P, int n, PT *V) {
int a, b, c;
int dnA, dnC;
if (above(P, V[n - 1], V[0]) && !below(P, V[1], V[0]))
 return 0;
for (a = 0, b = n;;) {
 c = (a + b) / 2;
dnC = below(P, V[c + 1], V[c]);
 if (above(P, V[c - 1], V[c]) && !dnC)
```

```
return c;
dnA = below(P, V[a + 1], V[a]);
if (dnA) {
   if (!dnC) {
      b = c;
   } else {
      if (below(P, V[a], V[c]))
      b = c;
      else
      a = c;
   }
} else {
   if (dnC) {
      a = c;
   } else {
   if (above(P, V[a], V[c]))
      b = c;
   else
      a = c;
}
}
```

6.14 Convex Hulls Tangent

```
// RLtangent_PolyPolyC(): get the RL tangent between
    two convex polygons
    Input: m = number of vertices in polygon 1
        V = array of vertices for convex polygon 1 with
     V[m]=V[0]
        n = number of vertices in polygon 2
        W = array of vertices for convex polygon 2 with
     W[n]=W[0]
    Output: *t1 = index of tangent point V[t1] for
    polygon 1
        *t2 = index of tangent point W[t2] for polygon
void RLtangent_PolyPolyC(int m, PT *V, int n, PT *W,
    int *t1, int *t2) {
 int ix1, ix2; // search indices for polygons 1 and 2
 // first get the initial vertex on each polygon
 ix1 = Rtangent_PointPolyC(W[0], m, V); // right
    tangent from W[0] to V
 ix2 = Ltangent_PointPolyC(V[ix1], n, W); // left
    tangent from V[ix1] to W
 // ping-pong linear search until it stabilizes
 int done = false; // flag when done
 while (done == false) {
  done = true; // assume done until..
  while (ori(W[ix2], V[ix1], V[ix1 + 1]) <= 0) {</pre>
   ++ix1; // get Rtangent from W[ix2] to V
  while (ori(V[ix1], W[ix2], W[ix2 - 1]) >= 0) {
            // get Ltangent from V[ix1] to W
   done = false; // not done if had to adjust this
 *t1 = ix1;
 *t2 = ix2:
```

6.15 Tangent line of Two Circle

```
vector<Line>
tanline(const Circle &c1, const Circle &c2, int sign1){
 // sign1 = 1 for outer tang, -1 for inter tang
 vector<Line> ret:
 if (norm(c1.o - c2.o) < eps) return ret;</pre>
 11f d = abs(c1.o - c2.o);
 PTF v = (c2.o - c1.o) / d;
 11f c = (c1.r - sign1 * c2.r) / d;
 if (c * c > 1) return ret;
 11f h = sqrt(max<11f>(0, 1 - c * c));
 for (int sign2 : {1, -1}) {
  PTF n = c * v + sign2 * h * rot90(v);
  PTF p1 = c1.o + n * c1.r;
  PTF p2 = c2.0 + n * (c2.r * sign1);
  if (norm(p2 - p1) < eps)
   p2 = p1 + rot90(c2.o - c1.o);
  ret.push_back({p1, p2});
```

```
return tree+M;
return ret;
}
                                                                  int touch(Node* r, int x, int y, LL d2){
                                                                   LL dis = sqrt(d2)+1;
6.16 Minimum Covering Circle
                                                                   if (x<r->x1-dis || x>r->x2+dis ||
template<typename P>
                                                                     y<r->y1-dis || y>r->y2+dis)
Circle getCircum(const P &a, const P &b, const P &c){
                                                                    return 0;
 Real a1 = a.x-b.x, b1 = a.y-b.y;
                                                                   return 1;
 Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
 Real a2 = a.x-c.x, b2 = a.y-c.y;
                                                                  void nearest(Node* r,int x,int y,int &mID,LL &md2) {
 Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
                                                                   if (!r || !touch(r, x, y, md2)) return;
                                                                   LL d2 = dis2(r->x, r->y, x, y);
 cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
                                                                   if (d2 < md2 | | (d2 == md2 && mID < r->id)) {
 cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2)
                                                                    mID = r->id;
                                                                    md2 = d2;
 cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
 return cc;
                                                                   // search order depends on split dim
                                                                   if ((r->f == 0 && x < r->x) ||
                                                                     (r->f == 1 \&\& y < r->y)) {
template<typename P>
                                                                    nearest(r->L, x, y, mID, md2);
nearest(r->R, x, y, mID, md2);
Circle MinCircleCover(const vector<P>& pts){
 random_shuffle(pts.begin(), pts.end());
 Circle c = { pts[0], 0 };
                                                                   } else {
 for(int i=0;i<(int)pts.size();i++){</pre>
                                                                    nearest(r->R, x, y, mID, md2);
 if (dist(pts[i], c.o) <= c.r) continue;</pre>
                                                                    nearest(r->L, x, y, mID, md2);
  c = { pts[i], 0 };
  for (int j = 0; j < i; j++) {
  if(dist(pts[j], c.o) <= c.r) continue;</pre>
                                                                  int query(int x, int y) {
   c.o = (pts[i] + pts[j]) / 2;
                                                                   int id = 1029384756;
   c.r = dist(pts[i], c.o);
                                                                   LL d2 = 102938475612345678LL;
   for (int k = 0; k < j; k++) {
                                                                   nearest(root, x, y, id, d2);
   if (dist(pts[k], c.o) <= c.r) continue;</pre>
                                                                   return id;
    c = getCircum(pts[i], pts[j], pts[k]);
                                                                 } tree;
  }
                                                                        Rotating Sweep Line
                                                                 6.18
                                                                 void rotatingSweepLine(pair<int, int> a[], int n) {
 return c;
                                                                  vector<pair<int, int>> 1;
                                                                  1.reserve(n * (n - 1) / 2)
       KDTree (Nearest Point)
                                                                  for (int i = 0; i < n; ++i)
for (int j = i + 1; j < n; ++j)
const int MXN = 100005;
struct KDTree {
                                                                    1.emplace_back(i, j);
                                                                  sort(1.begin(), 1.end(), [&a](auto &u, auto &v){
 struct Node {
                                                                   1ld udx = a[u.first].first - a[u.second].first;
  int x,y,x1,y1,x2,y2;
 int id,f;
Node *L, *R;
                                                                   11d udy = a[u.first].second - a[u.second].second;
                                                                   lld vdx = a[v.first].first - a[v.second].first;
lld vdy = a[v.first].second - a[v.second].second;
 } tree[MXN], *root;
                                                                   if (udx == 0 or vdx == 0) return not udx == 0;
                                                                   int s = sgn(udx * vdx);
LL dis2(int x1, int y1, int x2, int y2) {
LL dx = x1-x2, dy = y1-y2;
                                                                   return udy * vdx * s < vdy * udx * s;
 return dx*dx+dy*dy;
                                                                  }):
                                                                  vector<int> idx(n), p(n);
                                                                  iota(idx.begin(), idx.end(), 0);
sort(idx.begin(), idx.end(), [&a](int i, int j){
 static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
 static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 void init(vector<pair<int,int>> ip) {
                                                                   return a[i] < a[j]; });
                                                                  for (int i = 0; i < n; ++i) p[idx[i]] = i;
for (auto [i, j]: 1) {</pre>
 n = ip.size();
  for (int i=0; i<n; i++) {</pre>
   tree[i].id = i;
                                                                   // do here
   tree[i].x = ip[i].first;
                                                                   swap(p[i], p[j]);
   tree[i].y = ip[i].second;
                                                                   idx[p[i]] = i, idx[p[j]] = j;
  root = build_tree(0, n-1, 0);
                                                                 6.19
                                                                        Circle Cover
 Node* build_tree(int L, int R, int d) {
                                                                 const int N = 1021;
  if (L>R) return nullptr
  int M = (L+R)/2; tree[M].f = d%2;
                                                                 struct CircleCover {
  nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
                                                                  int C;
  tree[M].x1 = tree[M].x2 = tree[M].x;
                                                                  Cir c[N]
                                                                  bool g[N][N], overlap[N][N];
// Area[i] : area covered by at least i circles
  tree[M].y1 = tree[M].y2 = tree[M].y;
  tree[M].L = build_tree(L, M-1, d+1);
                                                                  double Area[ N ];
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
                                                                  void init(int _C){ C = _C;}
                                                                  struct Teve {
   tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
                                                                   PTF p; double ang; int add;
   tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
                                                                   Teve() {}
                                                                   Teve(PTF _a, double _b, int _c):p(_a), ang(_b), add(
  tree[M].R = build_tree(M+1, R, d+1);
  if (tree[M].R) {
                                                                   bool operator<(const Teve &a)const
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
                                                                   {return ang < a.ang;}
   tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
                                                                  }eve[N * 2];
   tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
                                                                  // strict: x = 0, otherwise x = -1
                                                                  bool disjuct(Cir &a, Cir &b, int x)
```

{return sign(abs(a.0 - b.0) - a.R - b.R) > x;}

```
bool contain(Cir &a, Cir &b, int x)
{return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
bool contain(int i, int j) {
 /* c[j] is non-strictly in c[i]. */
 return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c
   [j].R) == 0 && i < j)) && contain(c[i], c[j], -1);
void solve(){
 fill_n(Area, C + 2, 0);
 for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)
   overlap[i][j] = contain(i, j);
 for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j</pre>
   g[i][j] = !(overlap[i][j] \mid\mid overlap[j][i] \mid\mid
      disjuct(c[i], c[j], -1));
 for(int i = 0; i < C; ++i){
  int E = 0, cnt = 1;
  for(int j = 0; j < C; ++j)
   if(j != i && overlap[j][i])
    ++cnt;
  for(int j = 0; j < C; ++j)
   if(i != j && g[i][j]) {
    auto IP = intersectPoint(c[i], c[j]);
    PTF aa = IP[0], bb = IP[1];
    llf A = arg(aa-c[i].0), B = arg(bb-c[i].0);
eve[E++] = Teve(bb,B,1), eve[E++]=Teve(aa,A,-1);
    if(B > A) ++cnt;
  if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
   sort(eve, eve + E);
   eve[E] = eve[0];
   for(int j = 0; j < E; ++j){
    cnt += eve[j].add;
    Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
    double theta = eve[j + 1].ang - eve[j].ang;
    if (theta < 0) theta += 2. * pi;</pre>
    Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
 }
}
```

7 Stringology

7.1 Suffix Array

```
namespace sfx {
bool _t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], _c[maxn * 2];
int x[maxn], _p[maxn], _q[maxn * 2];
// sa[i]: sa[i]-th suffix is the
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *a, int *c, int n, int z) {
memset(a, 0, sizeof(int) * n);
memcpy(x, c, sizeof(int) * z);
void induce(int *a,int *c,int *s,bool *t,int n,int z){
memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
  if (a[i] && !t[a[i] - 1])
   a[x[s[a[i] - 1]]++] = a[i] - 1;
 memcpy(x, c, sizeof(int) * z);
 for (int i = n - 1; i >= 0; --i)
if (a[i] && t[a[i] - 1])
   a[--x[s[a[i] - 1]]] = a[i] - 1;
void sais(int *s, int *a, int *p, int *q,
bool *t, int *c, int n, int z) {
bool uniq = t[n - 1] = true;
 int nn=0, nmxz=-1, *nsa = a+n, *ns=s+n, last=-1;
 memset(c, 0, sizeof(int) * z);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
 if (uniq) {
  for (int i = 0; i < n; ++i) a[--c[s[i]]] = i;
  return;
```

```
for (int i = n - 2; i >= 0; --i)
  t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 pre(a, c, n, z);
 for (int i = 1; i <= n - 1; ++i)
  if (t[i] && !t[i - 1])
   a[--x[s[i]]] = p[q[i] = nn++] = i;
 induce(a, c, s, t, n, z);
for (int i = 0; i < n; ++i)
  if (a[i] && t[a[i]] && !t[a[i] - 1]) {
  bool neq = last < 0 \mid | \setminus memcmp(s + a[i], s + last,
   (p[q[a[i]] + 1] - a[i]) * sizeof(int));
  ns[q[last = a[i]]] = nmxz += neq;
 sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
 pre(a, c, n, z);
for (int i = nn - 1; i >= 0; --i)
  a[--x[s[p[nsa[i]]]] = p[nsa[i]];
 induce(a, c, s, t, n, z);
void build(const string &s) {
 const int n = int(s.size());
 for (int i = 0; i < n; ++i) _s[i] = s[i];
 _s[n] = 0; // s shouldn't contain 0
 sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
for(int i = 0; i < n; ++i) rev[sa[i] = sa[i+1]] = i;</pre>
 int ind = hi[0] = 0;
 for (int i = 0; i < n; ++i) {</pre>
  if (!rev[i]) {
   ind = 0;
   continue;
  while (i + ind < n && \</pre>
   s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  hi[rev[i]] = ind ? ind-- : 0;
}}
7.2 Suffix Automaton
struct SuffixAutomaton {
 struct node -
  int ch[K], len, fail, cnt, indeg;
  node(int L = 0) : ch{}, len(L), fail(0), cnt(0),
    indeg(0) {}
 } st[N];
 int root, last, tot;
 void extend(int c) {
  int cur = ++tot;
  st[cur] = node(st[last].len + 1);
  while (last && !st[last].ch[c]) {
    st[last].ch[c] = cur;
    last = st[last].fail;
  if (!last) {
    st[cur].fail = root;
  } else {
    int q = st[last].ch[c];
    if (st[q].len == st[last].len + 1) {
      st[cur].fail = q;
    } else {
      int clone = ++tot;
      st[clone] = st[q];
      st[clone].len = st[last].len + 1;
      st[st[cur].fail = st[q].fail = clone].cnt = 0;
      while (last && st[last].ch[c] == q) {
         st[last].ch[c] = clone;
         last = st[last].fail;
    }
  st[last = cur].cnt += 1;
 void init(const char* s) {
  root = last = tot = 1;
  st[root] = node(0);
  for (char c; c = *s; ++s) extend(c - 'a');
 int q[N];
 void dp() {
  for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
```

int head = 0, tail = 0;

k1 = 1 < z1.size() ? z1[1] : 0;

k2 = n + 1 - 1 < z2.size() ? z2[n + 1 - 1] : 0;

```
for (int i = 1; i <= tot; i++)
                                                                    } else {
    if (st[i].indeg == 0) q[tail++] = i;
                                                                     1 = cntr - nu + 1;
  while (head != tail) {
                                                                     k1 = n + 1 - 1 < z3.size() ? z3[n + 1 - 1] : 0;
    int now = q[head++];
                                                                     k2 = 1 < z4.size() ? z4[1] : 0;
    if (int f = st[now].fail) {
                                                                    if (k1 + k2 >= 1)
      st[f].cnt += st[now].cnt;
      if (--st[f].indeg == 0) q[tail++] = f;
                                                                     reps.emplace_back(cntr, 1, k1, k2);
  }
                                                                       BWT
                                                                  7.7
 int run(const char* s) {
                                                                  struct BurrowsWheeler{
  int now = root;
  for (char c; c = *s; ++s) {
   if (!st[now].ch[c -= 'a']) return 0;
                                                                  #define SIGMA 26
                                                                  #define BASE 'a'
                                                                  vector<int> v[ SIGMA ];
    now = st[now].ch[c];
                                                                   void BWT(char* ori, char* res){
                                                                    // make ori -> ori + ori
  return st[now].cnt;
                                                                    // then build suffix array
} SAM;
                                                                   void iBWT(char* ori, char* res){
7.3 Z value
                                                                    for( int i = 0 ; i < SIGMA ; i ++ )</pre>
vector<int> Zalgo(const string &s) {
                                                                     v[ i ].clear();
                                                                    int len = strlen( ori );
for( int i = 0 ; i < len ; i ++ )</pre>
 vector<int> z(s.size(), s.size());
for (int i = 1, 1 = 0, r = 0; i < z[0]; ++i) {
  int j = clamp(r - i, 0, z[i - 1]);
  for (; i + j < z[0] and s[i + j] == s[j]; ++j);
  if (i + (z[i] = j) > r) r = i + z[1 = i];
                                                                     v[ ori[i] - BASE ].push_back( i );
                                                                    vector<int> a:
                                                                    for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
                                                                     for( auto j : v[ i ] ){
                                                                      a.push_back( j );
ori[ ptr ++ ] = BASE + i;
 return z:
                                                                    for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
7.4 Manacher
int z[maxn];
                                                                     ptr = a[ ptr ];
int manacher(const string& s) {
 string t =
                                                                    res[ len ] = 0;
 for(char c: s) t += c, t += '.';
 int 1 = 0, r = 0, ans = 0;
                                                                 } bwt;
 for (int i = 1; i < t.length(); ++i) {
  z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
                                                                  7.8 Palindromic Tree
  while (i - z[i] >= 0 \& i + z[i] < t.length()) {
                                                                  struct palindromic_tree{
   if(t[i - z[i]] == t[i + z[i]]) ++z[i];
                                                                   struct node{
   else break;
                                                                    int next[26],f,len;
                                                                    int cnt, num, st, ed;
  if (i + z[i] > r) r = i + z[i], l = i;
                                                                    node(int l=0):f(0),len(1),cnt(0),num(0) {
                                                                     memset(next, 0, sizeof(next)); }
 for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);</pre>
 return ans;
                                                                   vector<node> st:
                                                                   vector<char> s;
                                                                   int last,n;
7.5 Lexico Smallest Rotation
                                                                   void init(){
string mcp(string s) {
                                                                    st.clear();s.clear();last=1; n=0;
 int n = s.length();
                                                                    st.push_back(0);st.push_back(-1);
 s += s; int i = 0, j = 1;
                                                                    st[0].f=1;s.push_back(-1); }
 while (i < n && j < n) {</pre>
                                                                   int getFail(int x){
  int k = 0;
                                                                    while(s[n-st[x].len-1]!=s[n])x=st[x].f;
  while (k < n \&\& s[i + k] == s[j + k]) k++;
                                                                    return x;}
  ((s[i+k] \le s[j+k]) ? j : i) += k + 1;
                                                                   void add(int c){
  j += (i == j);
                                                                    s.push_back(c-='a'); ++n;
                                                                    int cur=getFail(last);
 return s.substr(i < n ? i : j, n);</pre>
                                                                    if(!st[cur].next[c]){
                                                                     int now=st.size();
                                                                     st.push_back(st[cur].len+2);
7.6 Main Lorentz
                                                                     st[now].f=st[getFail(st[cur].f)].next[c];
vector<tuple<tuple<size_t, size_t, int, int>>> reps;
                                                                     st[cur].next[c]=now;
void find_repetitions(const string &s, int shift = 0) {
                                                                     st[now].num=st[st[now].f].num+1;
 if (s.size() <= 1)
                                                                    last=st[cur].next[c];
 const size_t nu = s.size() / 2, nv = s.size() - nu;
                                                                    ++st[last].cnt;}
 string u = s.substr(0, nu), v = s.substr(nu);
                                                                   void dpcnt() {
 string ru(u.rbegin(), u.rend());
string rv(v.rbegin(), v.rend());
                                                                    for (int i=st.size()-1; i >= 0; i--)
                                                                     st[st[i].f].cnt += st[i].cnt;
 find_repetitions(u, shift);
 find_repetitions(v, shift + nu);
                                                                   int size(){ return st.size()-2;}
 auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u),
z3 = Zalgo(ru + '#' + rv), z4 = Zalgo(v);
                                                                  } pt;
                                                                  int main() {
 for (size_t cntr = 0; cntr < s.size(); cntr++) {</pre>
                                                                   string s; cin >> s; pt.init();
                                                                   for (int i=0; i<SZ(s); i++)
  size_t l; int k1, k2;
  if (cntr < nu) {</pre>
                                                                    int prvsz = pt.size(); pt.add(s[i]);
                                                                    if (prvsz != pt.size())
   1 = nu - cntr
```

int r = i, l = r - pt.st[pt.last].len + 1;

// pal @ [1,r]: s.substr(1, r-1+1)

8 Misc

8.1 Theorems

8.1.1 Sherman-Morrison formula

$$(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1 + v^{\mathsf{T}}A^{-1}u}$$

8.1.2 Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|{\rm det}(\tilde{L}_{rr})|.$

8.1.3 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on C

8.1.4 Cayley's Formula

- Given a degree sequence d_1,d_2,\dots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.5 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + d_2 + \ldots + d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

holds for all $1 \le k \le n$.

8.1.6 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex

8.1.7 Euler's planar graph formula

 $V - E + F = C + 1, E \le 3V - 6$ (?)

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2)$, find maximum $S\in I_1\cap I_2$. For each iteration, build the directed graph and find a shortest path from s to t.

- $s \to x : S \sqcup \{x\} \in I_1$
- $x \to t : S \sqcup \{x\} \in I_2$
- $y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1$ (y is in the unique circuit of $S \sqcup \{x\}$)
- $x \to y: S \setminus \{y\} \sqcup \{x\} \in I_2$ (y is in the unique circuit of $S \sqcup \{x\}$)

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|.$ In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

8.2 Bitset LCS

```
scanf("%d%d", &n, &m), u = n / 64 + 1;
for (int i = 1, c; i <= n; i++)
    scanf("%d", &c), p[c].set(i);
for (int i = 1, c; i <= m; i++) {
    scanf("%d", &c), (g = f) |= p[c];
    f.shift(), f.set(0);
    ((f = g - f) ^= g) &= g;
}
printf("%d\n", f.count());</pre>
```

8.3 Convex 1D/1D DP

```
struct segment {
 int i, 1, r
 segment() {}
 segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
 dp[0] = 0;
 deque<segment> dq; dq.push_back(segment(0, 1, n));
 for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i);</pre>
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
   f(i, dq.back().1) < f(dq.back().i, dq.back().1))
    dq.pop_back();
  if (dq.size())
   int d = 1 << 20, c = dq.back().1;
   while (d \gg 1) if (c + d \ll d, back().r)
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
  if (seg.1 <= n) dq.push_back(seg);</pre>
}
}
```

8.4 ConvexHull Optimization

```
mutable int64_t a, b, p;
 bool operator<(const L &r) const { return a < r.a; }</pre>
bool operator<(int64_t x) const { return p < x; }</pre>
struct DynamicHull : multiset<L, less<>> {
 static const int64_t kInf = 1e18;
 bool Isect(iterator x, iterator y)
 auto Div = [](int64_t a, int64_t b) {
  return a / b - ((a ^ b) < 0 && a % b); }</pre>
  if (y == end()) { x->p = kInf; return false; }
  if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(int64_t a, int64_t b) {
  auto z = insert(\{a, b, 0\}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
  while ((y = x) != begin() && (--x)->p >= y->p)
   Isect(x, erase(y));
 int64_t Query(int64_t x) {
 auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
```

8.5 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

8.6 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
   dfn[u]=low[u]=++dfs_idx;
   for(int i=0;i<(int)init_g[u].size();i++){
    int v=init_g[u][i];
   if(v==par[u]) continue;
   if(!dfn[v]){</pre>
```

```
par[v]=u;
                                                              tarjan(1);
   tarjan(v);
                                                              dfs(1,-1);
                                                             printf("%d\n", max(dp[1][0], dp[1][1]));
  low[u]=min(low[u],low[v]);
   if(dfn[u]<low[v]){</pre>
                                                              return 0;
   g[u].push_back(v);
    g[v].push_back(u);
                                                            8.7 Tree Knapsack
  }else{
                                                            int dp[N][K]; PII obj[N];
                                                            vector<int> G[N];
  low[u]=min(low[u],dfn[v]);
   if(dfn[v]<dfn[u]){</pre>
                                                            void dfs(int u, int mx){
    int temp_v=u;
                                                              for(int s: G[u]) {
    bcc_id++;
                                                               if(mx < obj[s].first) continue;</pre>
   while(temp_v!=v){
                                                               for(int i=0;i<=mx-obj[s].FF;i++)</pre>
     g[bcc_id+n].push_back(temp_v);
                                                               dp[s][i] = dp[u][i];
     g[temp_v].push_back(bcc_id+n);
                                                               dfs(s, mx - obj[s].first);
     temp_v=par[temp_v];
                                                               for(int i=obj[s].FF;i<=mx;i++)</pre>
                                                                dp[u][i] = max(dp[u][i],
   g[bcc_id+n].push_back(v);
                                                                 dp[s][i - obj[s].FF] + obj[s].SS);
    g[v].push_back(bcc_id+n);
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
                                                            }
                                                            8.8 N Queens Problem
                                                            vector< int > solve( int n ) {
                                                              // no solution when n=2, 3
int dp[maxn][2],min_dp[2][2],tmp[2][2],tp[2];
                                                              vector< int > ret;
void dfs(int u,int fa){
                                                              if ( n % 6 == 2 ) {
if(u<=n){
                                                               for ( int i = 2 ; i <= n ; i += 2 )
 for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                ret.push_back( i );
   int v=g[u][i];
                                                               ret.push_back( 3 ); ret.push_back( 1 );
   if(v==fa) continue;
                                                               for ( int i = 7 ; i <= n ; i += 2 )
   dfs(v,u);
                                                                ret.push_back( i );
  memset(tp,0x8f,sizeof tp);
                                                               ret.push_back( 5 )
   if(v<=n){
                                                              } else if ( n % 6 == 3 ) {
    tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
                                                               for ( int i = 4 ; i <= n ; i += 2 )
    tp[1]=max(
                                                                ret.push_back( i );
     dp[u][0]+dp[v][0]+1
                                                               ret.push_back( 2 );
     dp[u][1]+max(dp[v][0],dp[v][1])
                                                               for ( int i = 5 ; i <= n ; i += 2 )
                                                                ret.push_back( i );
   }else{
                                                               ret.push_back( 1 ); ret.push_back( 3 );
    tp[0]=dp[u][0]+dp[v][0];
                                                              } else {
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
                                                               for ( int i = 2 ; i <= n ; i += 2 )
                                                               ret.push_back( i );
   dp[u][0]=tp[0],dp[u][1]=tp[1];
                                                               for ( int i = 1 ; i <= n ; i += 2 )
                                                                ret.push_back( i );
}else{
 for(int i=0;i<(int)g[u].size();i++){</pre>
                                                              return ret;
  int v=g[u][i];
  if(v==fa) continue;
  dfs(v,u);
                                                             8.9 Binary Search On Fraction
                                                            struct Q {
 min_dp[0][0]=0;
                                                             11 p. a:
 min_dp[1][1]=1;
                                                             Q go(Q b, 11 d) { return \{p + b.p*d, q + b.q*d\}; \}
 min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f3f;
  for(int i=0;i<(int)g[u].size();i++){</pre>
                                                            bool pred(Q);
   int v=g[u][i];
                                                            // returns smallest p/q in [lo, hi] such that
  if(v==fa) continue;
                                                            // pred(p/q) is true, and 0 <= p,q <= N
  memset(tmp,0x8f,sizeof tmp);
                                                            Q frac_bs(11 N) {
   tmp[0][0]=max(
                                                              Q lo{0, 1}, hi{1, 0};
   min_dp[0][0]+max(dp[v][0],dp[v][1]),
                                                              if (pred(lo)) return lo;
    min_dp[0][1]+dp[v][0]
                                                              assert(pred(hi));
                                                              bool dir = 1, L = 1, H = 1;
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
                                                              for (; L || H; dir = !dir) {
   tmp[1][0]=max(
                                                              ll len = 0, step = 1
   min_dp[1][0]+max(dp[v][0],dp[v][1]),
                                                               for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
   min_dp[1][1]+dp[v][0]
                                                                if (Q mid = hi.go(lo, len + step)
                                                                  mid.p > N || mid.q > N || dir ^ pred(mid))
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
                                                                 t++;
  memcpy(min_dp,tmp,sizeof tmp);
                                                                else len += step;
                                                               swap(lo, hi = hi.go(lo, len));
 dp[u][1]=max(min_dp[0][1], min_dp[1][0]);
                                                               (dir ? L : H) = !!len;
  dp[u][0]=min_dp[0][0];
                                                              return dir ? hi : lo;
int main(){
int m,a,b;
scanf("%d%d",&n,&m);
for(int i=0;i<m;i++){
  scanf("%d%d",&a,&b);</pre>
 init_g[a].push_back(b);
 init_g[b].push_back(a);
par[1]=-1;
```