Contents

1	Dasi	io.												1
1	Basi													
	1.1	vimrc												
	1.2	Debug Macro			 	٠				٠				
	1.3	Increase Stack			 									. 1
	1.4	Pragma Optimization			 									. 2
	1.5	IO Optimization			 				 					. 2
		· · · ·												
2	Date	a Structure												2
	2.1	Dark Magic												. 2
	2.2	Link-Cut Tree												
	2.3	LiChao Segment Tree												
		•												
	2.4	Treap												
	2.5	Linear Basis												
	2.6	Binary Search On Segment Tree			 					٠	٠	•		. 3
3	Gra	ph												4
	3.1	2-SAT (SCC)			 									. 4
	3.2	BCC			 									. 4
	3.3	Centroid Decomposition			 									. 4
	3.4	Directed Minimum Spanning Tree .												
	3.5	Dominator Tree												
	3.6	Edge Coloring												
	3.7	•												
		Lowbit Decomposition												
	3.8	Manhattan Minimum Spanning Tree												
	3.9	MaxClique												
		MaxCliqueDyn												
	3.11	Minimum Mean Cycle			 									
	3.12	Minimum Steiner Tree			 									. 7
		Mo's Algorithm on Tree												
		Virtual Tree												
	•		•	•	•	•	•	•	 •	•	•	•	•	
4	Mate	ching & Flow												8
•	4.1	Bipartite Matching												
	4.1	Dijkstra Cost Flow												
		•												
	4.3	Dinic												
	4.4	Flow Models												
	4.5	General Graph Matching			 	٠				٠				. 9
	4.6	Global Min-Cut			 									. 10
	4.7	GomoryHu Tree			 									. 10
	4.8	Kuhn Munkres			 									. 10
	4.9	Minimum Cost Circulation												
	4 10	Minimum Cost Maximum Flow												
		Maximum Weight Graph Matching .												
	4.11	Maximom weight Graph Matching .	•	•	 ٠.	•		•	 •	٠	•	•		. "
5	Mat	h												13
5	Mat													13 13
5	Mat 5.1	Strling Number												. 13
5		Strling Number			 									. 13 . 13
5	5.1	Strling Number 5.1.1 First Kind 5.1.2 Second Kind			 		 		 					. 13 . 13 . 13
5	5.1 5.2	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd			 		 		 				 	. 13 . 13 . 13
5	5.1 5.2 5.3	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey			 				 				 	. 13 . 13 . 13 . 13
5	5.1 5.2 5.3 5.4	Strling Number			 				 				 	. 13 . 13 . 13 . 13 . 13
5	5.1 5.2 5.3	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey			 				 				 	. 13 . 13 . 13 . 13 . 13
5	5.1 5.2 5.3 5.4	Strling Number			 				 				 	. 13 . 13 . 13 . 13 . 13 . 13
5	5.1 5.2 5.3 5.4 5.5	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn			 				 					. 13 . 13 . 13 . 13 . 13 . 13 . 13
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog			 									. 13 . 13 . 13 . 13 . 13 . 13 . 13
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 13
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 14 . 14
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 14 . 14
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 14 . 14 . 14
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 14 . 14 . 14 . 14
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Miller Rabin												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 14 . 14 . 14 . 14 . 15 . 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 14 . 14 . 14 . 15 . 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 14 . 14 . 14 . 15 . 15 . 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve)												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 14 . 14 . 14 . 14 . 15 . 15 . 15 . 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 14 . 14 . 14 . 14 . 15 . 15 . 15 . 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve)												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 14 . 14 . 14 . 14 . 15 . 15 . 15 . 15 . 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 14 . 14 . 14 . 15 . 15 . 15 . 15 . 15 . 15
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue												. 13 . 133 . 133 . 133 . 133 . 134 . 134 . 144 . 144 . 155 . 155 . 155 . 156 . 156 . 156 . 157 . 156 . 157 . 157 . 158 . 158
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex												. 13 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 145 . 155 . 155 . 155 . 156 . 166 . 177 . 177
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue												. 13 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 145 . 155 . 155 . 155 . 156 . 166 . 177 . 177
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.10 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.19 5.19 5.19 5.19 5.19 5.19 5.19	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction												. 13 . 133 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 155 . 155 . 155 . 156 . 156 . 157 . 177 . 177
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 Geo	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 13
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 5.21 Geo 6.1	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction metry Basic Geometry												. 13 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 155 . 155 . 155 . 156 . 157 . 177 . 177 . 177
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection												. 13 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 155 . 155 . 155 . 156 . 166 . 177 . 177 . 177 . 177
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 Geo 6.1 6.2 6.3	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull												. 13 . 133 . 133 . 133 . 133 . 134 . 135 . 144 . 144 . 155 . 155 . 155 . 156 . 166 . 167 . 177 . 177 . 177 . 178 . 178 . 188 . 188 . 188 . 198 . 198
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull												. 13 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 155 . 155 . 155 . 155 . 157 . 177 . 177 . 177 . 188 . 188 . 188 . 188
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair												. 13 . 133 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 155 . 155 . 155 . 156 . 157 . 177 . 177 . 177 . 177 . 188 . 188 . 188 . 188 . 188 . 188
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.14 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair kb Closest Pair (3D ver.)												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 13
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 13
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.14 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair kb Closest Pair (3D ver.)												. 13 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 155 . 155 . 155 . 156 . 166 . 166 . 177 . 177 . 177 . 177 . 188 . 188 . 188 . 188 . 188 . 188 . 188 . 188 . 189 . 189
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.12 5.14 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 3D Convex Hull 3D Closest Pair (3D ver.) Simulated Annealing												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 13
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.15 5.16 5.17 5.18 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijin DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum												. 13 . 133 . 133 . 133 . 133 . 134 . 134 . 144 . 144 . 155 . 155 . 155 . 155 . 157 . 177 . 177 . 177 . 178 . 188 . 188 . 188 . 199 . 199 . 199 . 199 . 199 . 199
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.6 6.7 6.6 6.7 6.7 6.7 6.7 6.7	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair kD Closest Pair Simulated Annealing Half Plane Intersection Minkowski Sum Circle Class												. 13 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 155 . 155 . 155 . 155 . 157 . 177 . 177 . 177 . 178 . 188 . 188 . 188 . 199 . 199 . 199 . 199 . 199 . 199 . 199
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.14 5.16 5.17 5.18 5.19 5.21 Geo 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.9 6.0 6.0 6.0 6.0 6.0 6.0 6.0 6.0 6.0 6.0	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Circle Class Intersection of line and Circle												. 13 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 155 . 155 . 155 . 156 . 157 . 177 . 177 . 177 . 177 . 188 . 188 . 188 . 189 . 199 . 199 . 199 . 199 . 199 . 199 . 199 . 199 . 199
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.14 5.15 5.16 5.17 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 6.12 6.11 6.12 6.12 6.13 6.14 6.14 6.15 6.15 6.15 6.15 6.15 6.15 6.15 6.15	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Farthest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Circle Class Intersection of Polygon and Circle Intersection of Polygon and Circle												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 13
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.10 5.11 5.12 5.14 5.15 5.16 5.17 5.19 5.20 6.1 6.2 6.3 6.4 6.6 6.7 6.8 6.9 6.10 6.11 6.12 6.13	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 3D Farthest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Circle Class Intersection of Polygon and Circle Intersection of Polygon and Circle Point & Hulls Tangent												. 13 . 13 . 13 . 13 . 13 . 13 . 13 . 13
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.19 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.10 6.11 6.12 6.13 6.14 6.13 6.14 6.14 6.14 6.14 6.14 6.14 6.14 6.14	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 2D Farthest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Circle Class Intersection of line and Circle Intersection of Polygon and Circle Point & Hulls Tangent Convex Hulls Tangent												. 13 . 133 . 133 . 133 . 133 . 134 . 135 . 144 . 144 . 155 . 155 . 155 . 156 . 166 . 167 . 177 . 177 . 177 . 179 . 188 . 188 . 188 . 188 . 189 . 199 . 199
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.10 6.11 6.12 6.13 6.14 6.14 6.15 6.16 6.16 6.16 6.16 6.16 6.16 6.16	Strling Number 5.1.1 First Kind 5.1.2 Second Kind 3.1.2 Second Kind 3.1.4 Second Kind 3.1.5 Second Kind 3.1.5 Second Kind 3.1.6 Second Kind 3.1.7 Second Kind 3.1 Second Kind 3.												. 13 . 133 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 155 . 155 . 155 . 155 . 157 . 177 . 177 . 177 . 177 . 188 . 188 . 188 . 199 . 199
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.14 5.15 5.15 5.16 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.1 6.1 6.1 6.1 6.1 6.1 6.1 6.1 6.1 6.1	Strling Number 5.1.1 First Kind 5.1.2 Second Kind ax+by=gcd Berlekamp Massey Charateristic Polynomial Chinese Remainder De-Bruijn DiscreteLog Extended Euler Extended FloorSum Fast Fourier Transform FloorSum FWT Miller Rabin NTT Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair kD Closest Pair kD Clos												. 13 . 133 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 155 . 155 . 155 . 155 . 157 . 177 . 177 . 177 . 177 . 178 . 188 . 188 . 188 . 199 . 199 . 199 . 199 . 199 . 199 . 199 . 200 . 200
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.11 5.12 5.13 5.14 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.10 6.11 6.12 6.13 6.14 6.15 6.16 6.16 6.17 6.16 6.17 6.17 6.17 6.17	Strling Number 5.1.1 First Kind 5.1.2 Second Kind 3.1.2 Second Kind 3.1.4 Second Kind 3.1.5 Second Kind 3.1.5 Second Kind 3.1.6 Second Kind 3.1.7 Second Kind 3.1 Second Kind 3.												. 13 . 133 . 133 . 133 . 133 . 133 . 134 . 144 . 144 . 144 . 155 . 155 . 155 . 156 . 166 . 167 . 177 . 177 . 177 . 177 . 178 . 188 . 188 . 188 . 189 . 199 . 199

Strir	agalogu	22
	5 55	22 22
		22
		22
		23
		23
		23
		23
		23 23
7.0	rumuroniic nee	23
Misc		23
8.1	Theorems	23
	8.1.1 Sherman-Morrison formula	23
	8.1.2 Kirchhoff's Theorem	23
	8.1.3 Tutte's Matrix	23
	8.1.4 Cayley's Formula	23
	8.1.5 Erdős-Gallai theorem	24
	8.1.6 Havel-Hakimi algorithm	24
		24
	8.1.8 Pick's theorem	24
	8.1.9 Matroid Intersection	24
8.2	Bitset LCS	24
8.3	Convex 1D/1D DP	24
8.4	ConvexHull Optimization	24
8.5	Josephus Problem	24
8.6		24
8.7		25
8.8	N Queens Problem	25
8.9	·-	25
	-	
F	Rasic	
	J431C	
	7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8 Misc 8.1 8.2 8.3 8.4 8.5 8.6 8.7 8.8 8.9	7.1 Suffix Array 7.2 Suffix Automaton 7.3 Z value 7.4 Manacher 7.5 Lexico Smallest Rotation 7.6 Main Lorentz 7.7 BWT 7.8 Palindromic Tree Misc 8.1 Theorems 8.1.1 Sherman-Morrison formula 8.1.2 Kirchhoff's Theorem 8.1.3 Tutte's Matrix 8.1.4 Cayley's Formula 8.1.5 Erdős-Gallai theorem 8.1.6 Havel-Hakimi algorithm 8.1.7 Euler's planar graph formula 8.1.8 Pick's theorem 8.1.9 Matroid Intersection 8.2 Bitset LCS 8.3 Convex 1D/1D DP 8.4 ConvexHull Optimization 8.5 Josephus Problem 8.6 Cactus Matching 8.7 Tree Knapsack 8.8 N Queens Problem

1.1 vimrc

```
se is nu ru et tgc sc hls cin cino+=j1 sw=4 sts=4 bs=2
    mouse=a encoding=utf-8 ls=2
syn on
colorscheme desert
filetype indent on
inoremap {<CR> {<CR>}<ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -std=c++17 -
    DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
    Wconversion -fsanitize=address,undefined -g && echo
    success<CR>
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -O2 -std=c++17 &&
    echo success<CR>
map <F10> <ESC>:!./"%<" <CR>
```

1.2 Debug Macro

```
#ifdef KISEKI
#define safe cerr<<__PRETTY_FUNCTION__\
    <<" line "<<__LINE__<<" safe\n"</pre>
#define debug(a...) qwerty(#a, a)
#define orange(a...) dvorak(#a, a)
using std::cerr;
template <typename ...T>
void qwerty(const char *s, T ...a) {
  cerr << "\e[1;32m(" << s << ") = (";</pre>
  int cnt = sizeof...(T);
  (..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
template <typename Iter>
void dvorak(const char *s, Iter L, Iter R) {
  cerr << "\e[1;32m[ " << s << " ] = [ ";</pre>
  for (int f = 0; L != R; ++L)
  cerr << (f++ ? ", " : "") << *L;
cerr << " ]\e[0m\n";</pre>
#else
#define safe ((void)0)
#define debug(...) ((void)0)
#define orange(...) ((void)0)
#endif
```

1.3 Increase Stack

```
const int size = 256 << 20;
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));</pre>
```

1.4 Pragma Optimization

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,sse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8000)
;
```

1.5 IO Optimization

```
| static inline int gc() {
    constexpr int B = 1<<20;
    static char buf[B], *p, *q;
    if(p == q &&
        (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
        return EOF;
    return *p++;
    }
    template < typename T >
    static inline bool gn( T &x ) {
        int c = gc(); T sgn = 1; x = 0;
        while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
        if(c == '-') sgn = -1, c = gc();
        if(c == EOF) return false;
        while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
        return x *= sgn, true;
}
```

2 Data Structure

2.1 Dark Magic

2.2 Link-Cut Tree

```
template <typename Val> class LCT {
private:
struct node
  int pa, ch[2];
 bool rev;
 Val v, v_prod, v_rprod;
 node() \; : \; pa\{0\}, \; ch\{0, \; 0\}, \; rev\{false\}, \; v\{\}, \; v\_prod\{\}, \\
    v_rprod{} {};
};
vector<node> nodes;
 set<pair<int, int>> edges;
bool is_root(int u) const {
 const int p = nodes[u].pa;
 return nodes[p].ch[0] != u and nodes[p].ch[1] != u;
bool is_rch(int u) const {
 return (not is_root(u)) and nodes[nodes[u].pa].ch[1]
    == u;
void down(int u) {
 if (auto &cnode = nodes[u]; cnode.rev) {
   if (cnode.ch[0]) set_rev(cnode.ch[0]);
   if (cnode.ch[1]) set_rev(cnode.ch[1]);
   cnode.rev = false;
 void up(int u) {
 auto &cnode = nodes[u];
 cnode.v_prod =
  nodes[cnode.ch[0]].v_prod * cnode.v * nodes[cnode.ch
    [1]].v_prod;
  cnode.v_rprod =
   nodes[cnode.ch[1]].v_rprod * cnode.v * nodes[cnode.
    ch[0]].v_rprod;
```

```
void set_rev(int u) {
  swap(nodes[u].ch[0], nodes[u].ch[1]);
  swap(nodes[u].v_prod, nodes[u].v_rprod);
 nodes[u].rev ^= 1;
 void rotate(int u) {
 int f = nodes[u].pa, g = nodes[f].pa, l = is_rch(u);
if (nodes[u].ch[l ^ 1])
  nodes[nodes[u].ch[1 ^ 1]].pa = f;
  if (not is_root(f))
   nodes[g].ch[is_rch(f)] = u;
  nodes[f].ch[1] = nodes[u].ch[1 ^ 1];
  nodes[u].ch[1 ^ 1] = f
  nodes[u].pa = g, nodes[f].pa = u;
 up(f);
 void splay(int u) {
  vector<int> stk = {u};
  while (not is_root(stk.back()))
  stk.push_back(nodes[stk.back()].pa);
  for (; not stk.empty(); stk.pop_back())
   down(stk.back())
  for(int f=nodes[u].pa;!is_root(u);f=nodes[u].pa){
   if(!is_root(f))rotate(is_rch(u)==is_rch(f)?f:u);
   rotate(u);
  }
 up(u);
 void access(int u) {
  int last = 0;
  for (int last = 0; u; last = u, u = nodes[u].pa) {
   splay(u);
   nodes[u].ch[1] = last;
   up(u);
 int find_root(int u) {
  access(u); splay(u);
  int la = 0:
  for (; u; la = u, u = nodes[u].ch[0]) down(u);
  return la;
 void change_root(int u) {
 access(u); splay(u); set_rev(u);
 void link(int x, int y) {
 change_root(y); nodes[y].pa = x;
 void split(int x, int y) {
 change_root(x); access(y); splay(y);
 void cut(int x, int y) {
 split(x, y)
  nodes[y].ch[0] = nodes[x].pa = 0;
 up(y);
public:
LCT(int n = 0) : nodes(n + 1) {}
 int add(const Val &v = {}) {
 nodes.push_back(v)
  return int(nodes.size()) - 2;
 int add(Val &&v) {
 nodes.emplace_back(move(v));
  return int(nodes.size()) - 2;
 void set_val(int u, const Val &v) {
 splay(++u); nodes[u].v = v; up(u);
 Val query(int x, int y) {
 split(++x, ++y)
  return nodes[y].v_prod;
 bool connected(int u, int v) { return find_root(++u)
    == find_root(++v); }
 void add_edge(int u, int v) {
 if (++u > ++v) swap(u, v)
  edges.emplace(u, v); link(u, v);
 void del_edge(int u, int v) {
 auto k = minmax(++u, ++v);
```

```
if (auto it = edges.find(k); it != edges.end()) {
                                                                } // sz(L) == k
   edges.erase(it); cut(u, v);
                                                                int getRank(node *o) { // 1-base
 }
                                                                 int r = sz(o->lc) + 1;
};
                                                                 for (;o->pa != nullptr; o = o->pa)
                                                                  if (o->pa->rc == o) r += sz(o->pa->lc) + 1;
2.3 LiChao Segment Tree
                                                                 return r:
struct L {
int m, k, id;
L() : id(-1) {}
L(int a, int b, int c) : m(a), k(b), id(c) {}
                                                                #undef sz
                                                               2.5 Linear Basis
 int at(int x) { return m * x + k; }
                                                               template <int BITS> struct Basis {
class LiChao {
                                                                array<pair<uint64_t, int>, BITS> b;
                                                                Basis() { b.fill({0, -1}); }
private:
                                                                void add(uint64_t x, int p) {
  for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
 int n; vector<L> nodes;
static int lc(int x) { return 2 * x + 1;
static int rc(int x) { return 2 * x + 2;
                                                                  if (b[i].first == 0) {
 void insert(int 1, int r, int id, L ln) {
                                                                  b[i] = \{x, p\};
  int m = (1 + r) >> 1;
                                                                   return;
  if (nodes[id].id == -1) {
                                                                  } else if (b[i].second > p) {
   nodes[id] = ln;
                                                                   swap(b[i].first, x), swap(b[i].second, p);
   return:
                                                                  x ^= b[i].first;
  bool atLeft = nodes[id].at(1) < ln.at(1);</pre>
                                                                 }
  if (nodes[id].at(m) < ln.at(m)) {</pre>
  atLeft ^= 1
                                                                bool ok(uint64_t x, int p) {
   swap(nodes[id], ln);
                                                                 for (int i = 0; i < BITS; ++i)</pre>
                                                                  if (((x >> i) \& 1) \text{ and } b[i].second < p)
                                                                   x ^= b[i].first;
  if (r - l == 1) return;
  if (atLeft) insert(l, m, lc(id), ln);
                                                                 return x == 0;
  else insert(m, r, rc(id), ln);
                                                              };
 int query(int 1, int r, int id, int x) {
                                                               2.6 Binary Search On Segment Tree
 int ret = 0, m = (1 + r) >> 1;
if (nodes[id].id != -1)
                                                               // find_first = x -> minimal x s.t. check( [a, x) )
                                                               // find_last = x \rightarrow maximal x s.t. check([x, b))
  ret = nodes[id].at(x);
  if (r - 1 == 1) return ret;
                                                               template <typename C>
  if (x < m) return max(ret, query(1, m, lc(id), x));
return max(ret, query(m, r, rc(id), x));</pre>
                                                               int find_first(int 1, const C &check) {
                                                                if (1 >= n) return n + 1;
                                                                1 += sz;
                                                                for (int i = height; i > 0; i--)
public:
                                                                 propagate(1 >> i);
 LiChao(int n_{-}) : n(n_{-}), nodes(n * 4) {}
                                                                Monoid sum = identity;
 void insert(L ln) { insert(0, n, 0, ln); }
                                                                do {
 int query(int x) { return query(0, n, 0, x); }
                                                                 while ((1 & 1) == 0) 1 >>= 1;
                                                                 if (check(f(sum, data[1]))) {
                                                                  while (1 < sz) {
2.4 Treap
                                                                   propagate(1);
namespace Treap{
                                                                   1 <<= 1;
 #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                                   auto nxt = f(sum, data[1]);
 struct node{
                                                                   if (not check(nxt)) {
                                                                    sum = nxt;
  int size:
  uint32_t pri;
                                                                    1++:
 node *lc, *rc, *pa;
                                                                   }
  node():size(0),pri(rand()),lc(0),rc(0),pa(0){}
  void pull() {
                                                                  return 1 + 1 - sz;
  size = 1; pa = nullptr;
   if ( lc ) { size += lc->size; lc->pa = this; }
                                                                 sum = f(sum, data[1++]);
   if ( rc ) { size += rc->size; rc->pa = this; }
                                                                } while ((1 & -1) != 1);
  }
                                                                return n + 1;
 node* merge( node* L, node* R ) {
                                                               template <typename C>
  if ( not L or not R ) return L ? L : R;
                                                               int find_last(int r, const C &check) {
  if ( L->pri > R->pri ) {
                                                                if (r <= 0) return -1;</pre>
  L->rc = merge( L->rc, R ); L->pull();
                                                                r += sz;
   return L;
                                                                for (int i = height; i > 0; i--)
  } else {
                                                                 propagate((r - 1) >> i);
   R->lc = merge( L, R->lc ); R->pull();
                                                                Monoid sum = identity;
                                                                do {
   return R;
  }
 }
                                                                 while (r > 1 and (r & 1)) r >>= 1;
 void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                                 if (check(f(data[r], sum))) {
 if ( not rt ) L = R = nullptr;
                                                                  while (r < sz) {</pre>
  else if( sz( rt->lc ) + 1 <= k ) {
                                                                   propagate(r);
                                                                   r = (r << 1) + 1;
                                                                   auto nxt = f(data[r], sum);
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
   L->pull();
                                                                   if (not check(nxt)) {
  } else {
                                                                    sum = nxt;
                                                                    r--;
   split_by_size( rt->lc, k, L, R->lc );
   R->pull();
```

if (dfn[v]) {

```
low[u] = min(low[u], dfn[v]);
   return r - sz;
                                                                 } else {
  sum = f(data[r], sum);
                                                                  ++ch, dfs(v, u);
 } while ((r & -r) != r);
                                                                  low[u] = min(low[u], low[v]);
                                                                  if (low[v] > dfn[u])
 return -1;
                                                                   bridge[t] = true
                                                                  if (low[v] >= dfn[u])
                                                                   ap[u] = true;
3
    Graph
3.1 2-SAT (SCC)
                                                                ap[u] &= (ch != 1 or u != f);
class TwoSat{
private:
                                                              public:
 int n;
                                                               void init(int n_) {
 vector<vector<int>> rG,G,sccs;
                                                                g.assign(n = n_, vector<pair<int, int>>());
 vector<int> ord,idx;
                                                                low.assign(n, ecnt = 0);
 vector<bool> vis,result;
                                                                dfn.assign(n, 0);
 void dfs(int u){
                                                                ap.assign(n, false);
  vis[u]=true
  for(int v:G[u])
                                                               void add_edge(int u, int v) {
   if(!vis[v]) dfs(v);
                                                                g[u].emplace_back(v, ecnt);
  ord.push_back(u);
                                                                g[v].emplace_back(u, ecnt++);
 void rdfs(int u){
                                                               void solve() {
 vis[u]=false;idx[u]=sccs.size()-1;
                                                                bridge.assign(ecnt, false);
  sccs.back().push_back(u);
                                                                for (int i = 0; i < n; ++i)
  if (not dfn[i]) dfs(i, i);</pre>
  for(int v:rG[u])
   if(vis[v])rdfs(v);
                                                               bool is_ap(int x) { return ap[x]; }
public:
                                                               bool is_bridge(int x) { return bridge[x]; }
 void init(int n_){
  G.clear();G.resize(n=n_);
  rG.clear();rG.resize(n);
                                                              3.3 Centroid Decomposition
  sccs.clear();ord.clear();
                                                              struct Centroid {
  idx.resize(n);result.resize(n);
                                                               vector<vector<int64_t>> Dist;
                                                               vector<int> Pa, Dep;
 void add_edge(int u,int v){
                                                               vector<int64_t> Sub, Sub2;
 G[u].push_back(v);rG[v].push_back(u);
                                                               vector<int> Cnt, Cnt2;
                                                               vector<int> vis, sz, mx, tmp
 void orr(int x,int y){
                                                               void DfsSz(int x) {
 if ((x^y)==1)return
                                                                vis[x] = true; sz[x] = 1; mx[x] = 0;
  add_edge(x^1,y); add_edge(y^1,x);
                                                                for (auto [u, w] : g[x]) {
  if (vis[u]) continue;
 bool solve(){
                                                                 DfsSz(u);
  vis.clear();vis.resize(n);
                                                                 sz[x] += sz[u]
  for(int i=0;i<n;++i)</pre>
                                                                 mx[x] = max(mx[x], sz[u]);
  if(not vis[i])dfs(i);
  reverse(ord.begin(),ord.end());
                                                                tmp.push_back(x);
  for (int u:ord){
   if(!vis[u])continue;
                                                               void DfsDist(int x, int64_t D = 0) {
   sccs.push_back(vector<int>());
                                                                Dist[x].push_back(D); vis[x] = true;
   rdfs(u);
                                                                for (auto [u, w] : g[x])
  if (not vis[u]) DfsDist(u, D + w);
  for(int i=0;i<n;i+=2)</pre>
   if(idx[i]==idx[i+1])
                                                               void DfsCen(int x, int D = 0, int p = -1) {
    return false;
                                                                tmp.clear(); DfsSz(x);
  vector<bool> c(sccs.size());
                                                                int M = tmp.size();
  for(size_t i=0;i<sccs.size();++i){</pre>
                                                                int C = -1;
   for(auto sij : sccs[i]){
                                                                for (int u : tmp) {
    result[sij]=c[i];
                                                                 if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
    c[idx[sij^1]]=!c[i];
                                                                 vis[u] = false;
   }
  }
                                                                DfsDist(C);
  return true;
                                                                for (int u : tmp) vis[u] = false;
Pa[C] = p; vis[C] = true; Dep[C] = D;
 bool get(int x){return result[x];}
                                                                for (auto [u, w] : g[C])
 int get_id(int x){return idx[x];}
                                                                 if (not vis[u]) DfsCen(u, D + 1, C);
 int count(){return sccs.size();}
} sat2;
                                                               Centroid(int N, vector<vector<pair<int,int>>> g)
                                                                : Sub(N), Sub2(N), Cnt(N), Cnt2(N), Dist(N),
3.2 BCC
                                                                Pa(N), Dep(N), vis(N), sz(N), mx(N)
class BCC {
                                                                { DfsCen(0); }
private:
                                                               void Mark(int v) {
 int n, ecnt;
                                                                int x = v, z = -1
                                                                for (int i = Dep[v]; i >= 0; --i) {
 vector<vector<pair<int, int>>> g;
 vector<int> dfn, low;
                                                                 Sub[x] += Dist[v][i]; Cnt[x]++;
 vector<bool> ap, bridge;
void dfs(int u, int f) {
                                                                 if (z != -1) +
                                                                  Sub2[z] += Dist[v][i];
  dfn[u] = low[u] = dfn[f] + 1;
                                                                  Cnt2[z]++;
  int ch = 0;
  for (auto [v, t] : g[u]) if (v != f) {
                                                                 z = x; x = Pa[x];
```

void merge(int x, int y) { fa[x] = y; }

```
int find(int x, int c = 0) {
                                                                   if (fa[x] == x) return c ? -1 : x;
 int64_t Query(int v) {
                                                                   int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
  int64_t res = 0;
  int x = v, z = -1
                                                                   if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
  for (int i = Dep[v]; i >= 0; --i) {
                                                                   fa[x] = p;
return c ? p : val[x];
   res += Sub[x] + 1LL * Cnt[x] * Dist[v][i];
   if (z != -1) res-=Sub2[z]+1LL*Cnt2[z]*Dist[v][i];
   z = x; x = Pa[x];
                                                                 vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
  return res:
 }
                                                                  // p[i] = -2 if i is unreachable from s
};
                                                                   dfs(s);
                                                                   for (int i = tk - 1; i >= 0; --i) {
  for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
3.4
      Directed Minimum Spanning Tree
                                                                    if (i) rdom[sdom[i]].push_back(i);
struct Edge { int u, v, w; };
struct DirectedMST { // find maximum
                                                                    for (int &u : rdom[i]) {
 int solve(vector<Edge> E, int root, int n) {
                                                                     int p = find(u);
                                                                     if (sdom[p] == i) dom[u] = i;
  int ans = 0:
                                                                     else dom[u] = p;
  while (true) {
   // find best in edge
   vector<int> in(n, -inf), prv(n, -1);
                                                                    if (i) merge(i, rp[i]);
   for (auto e : E)
                                                                   vector<int> p(n, -2); p[s] = -1;
for (int i = 1; i < tk; ++i)</pre>
    if (e.u != e.v && e.w > in[e.v]) {
     in[e.v] = e.w;
     prv[e.v] = e.u;
                                                                    if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
                                                                   for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
   in[root] = 0; prv[root] = -1;
                                                                   return p;
   for (int i = 0; i < n; i++)</pre>
    if (in[i] == -inf) return -inf;
                                                                  3.6 Edge Coloring
   // find cycle
                                                                  // max(d_u) + 1 edge coloring, time: O(NM)
   int tot = 0;
   vector < int > id(n, -1), vis(n, -1);
                                                                  int C[kN][kN], G[kN][kN]; // 1-based, G: ans
   for (int i = 0; i < n; i++) {
                                                                  void clear(int N) {
    ans += in[i];
                                                                   for (int i = 0; i <= N; i++)</pre>
                                                                    for (int j = 0; j <= \hat{N}; j++)
    for (int x = i; x != -1 && id[x] == -1; x = prv[x])
                                                                      C[i][j] = G[i][j] = 0;
     if (vis[x] == i) {
      for (int y = prv[x]; y != x; y = prv[y])
                                                                  void solve(vector<pair<int, int>> &E, int N) {
                                                                   int X[kN] = {}, a;
auto update = [&](int u) {
        id[y] = tot;
      id[x] = tot++;
      break;
                                                                    for (X[u] = 1; C[u][X[u]]; X[u]++);
     }
                                                                   auto color = [&](int u, int v, int c) {
     vis[x] = i;
                                                                    int p = G[u][v];
                                                                    G[u][v] = G[v][u] = c;
   if (!tot) return ans;
                                                                    C[u][c] = v, C[v][c] = u;
   for (int i = 0; i < n; i++)</pre>
                                                                    C[u][p] = C[v][p] = 0;
    if (id[i] == -1) id[i] = tot++;
                                                                    if(p) X[u] = X[v] = p
   for (auto &e : E) {
                                                                    else update(u), update(v);
    if (id[e.u] != id[e.v]) e.w -= in[e.v];
                                                                    return p:
    e.u = id[e.u], e.v = id[e.v];
                                                                   auto flip = [&](int u, int c1, int c2) {
                                                                    int p = C[u][c1];
   n = tot; root = id[root];
                                                                    swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
if (!C[u][c1]) X[u] = c1;
  }
} DMST:
                                                                    if (!C[u][c2]) X[u] = c2;
3.5 Dominator Tree
                                                                    return p:
namespace dominator {
                                                                   for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {
  auto [u, v] = E[t];</pre>
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
                                                                    int v0 = v, c = X[u], c0 = c, d;
 // vertices are numbered from 0 to n - 1
                                                                    vector<pair<int, int>> L; int vst[kN] = {};
 fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
                                                                    while (!G[u][v0]) {
 fill(fa, fa + n, -1); fill(val, val + n, -1);
                                                                     L.emplace_back(v, d = X[v]);
 fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
                                                                     if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
fill(dom, dom + n, -1); tk = 0;
for (int i = 0; i < n; ++i) {
                                                                     c = color(u, L[a].first, c);
else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
  g[i].clear(); r[i].clear(); rdom[i].clear();
                                                                       color(u, L[a].first, L[a].second);
                                                                     else if (vst[d]) break
                                                                     else vst[d] = 1, v = C[u][d];
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
                                                                    if (!G[u][v0]) {
 rev[dfn[x] = tk] = x;
                                                                     for (; v; v = flip(v, c, d), swap(c, d));
                                                                     if (C[u][c0]) { a = int(L.size()) - 1;
 fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
                                                                      while (--a >= 0 && L[a].second != c);
                                                                      for(;a>=0;a--)color(u,L[a].first,L[a].second);
  r[dfn[u]].push_back(dfn[x]);
                                                                     } else t--;
```

3.7 Lowbit Decomposition

```
class LBD {
 int timer, chains;
 vector<vector<int>> G;
 vector<int> t1, tr, chain, head, dep, pa;
 // chains : number of chain
 // tl, tr[u] : subtree interval in the seq. of u
 // head[i] : head of the chain i
 // chian[u] : chain id of the chain u is on
 void predfs(int u, int f) {
dep[u] = dep[pa[u] = f] + 1;
  for (int v : G[u]) if (v != f) {
   predfs(v, u);
   if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
    chain[u] = chain[v];
  if (chain[u] == 0) chain[u] = ++chains;
 }
 void dfschain(int u, int f) {
  tl[u] = timer++;
  if (head[chain[u]] == -1)
   head[chain[u]] = u;
  for (int v : G[u])
  if (v != f and chain[v] == chain[u])
    dfschain(v, u);
  for (int v : G[u])
  if (v != f and chain[v] != chain[u])
    dfschain(v, u);
  tr[u] = timer;
public:
 LBD(\textbf{int } n) \ : \ timer(0), \ chains(0), \ G(n), \ tl(n), \ tr(n),
 G[u].push_back(v); G[v].push_back(u);
 void decompose() { predfs(0, 0); dfschain(0, 0); }
 PII get_subtree(int u) { return {tl[u], tr[u]}; }
 vector<PII> get_path(int u, int v) {
  vector<PII> res;
  while (chain[u] != chain[v]) {
   if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
    swap(u, v)
   int s = head[chain[u]];
   res.emplace_back(tl[s], tl[u] + 1);
   u = pa[s];
  if (dep[u] < dep[v]) swap(u, v);</pre>
  res.emplace_back(tl[v], tl[u] + 1);
  return res;
};
```

Manhattan Minimum Spanning Tree

```
typedef Point<int> P;
vector<array<int, 3>> manhattanMST(vector<P> ps) {
vi id(sz(ps));
iota(all(id), 0);
vector<array<int, 3>> edges;
 rep(k, 0, 4) {
 sort(all(id), [&](int i, int j) {
  return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
 map<int, int> sweep;
  for (int i : id) {
  for (auto it = sweep.lower_bound(-ps[i].y);
      it != sweep.end(); sweep.erase(it++)) {
    int j = it->second;
    P d = ps[i] - ps[j];
    if (d.y > d.x) break;
    edges.push_back({d.y + d.x, i, j});
   sweep[-ps[i].y] = i;
  for (P &p : ps)
   if (k \& 1) p.x = -p.x;
   else swap(p.x, p.y);
return edges; // [{w, i, j}, ...]
```

3.9 MaxClique

```
// contain a self loop u to u, than u won't in clique
template < size_t MAXN >
class MaxClique{
private:
 using bits = bitset< MAXN >;
 bits popped, G[ MAXN ], ans
 size_t deg[ MAXN ], deo[ MAXN ], n;
 void sort_by_degree() {
  popped.reset();
  for ( size_t i = 0 ; i < n ; ++ i )
    deg[ i ] = G[ i ].count();
  for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;</pre>
    for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )</pre>
        mi = deg[id = j];
    popped[ deo[ i ] = id ] = 1;
    for( size_t u = G[ i ]._Find_first() ;
  u < n ; u = G[ i ]._Find_next( u ) )</pre>
       -- deg[ u ];
  }
 void BK( bits R, bits P, bits X ) {
  if (R.count()+P.count() <= ans.count()) return;</pre>
  if ( not P.count() and not X.count() )
   if ( R.count() > ans.count() ) ans = R;
   return:
  /* greedily chosse max degree as pivot
  bits cur = P | X; size_t pivot = 0, sz = 0;
  for ( size_t u = cur._Find_first() ;
   u < n ; u = cur._Find_next( u ) )</pre>
  if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
cur = P & ( ~G[ pivot ] );
  */ // or simply choose first
  bits cur = P & (~G[ ( P | X )._Find_first() ]);
  for ( size_t u = cur._Find_first()
   u < n ; u = cur._Find_next( u ) ) {
   if ( R[ u ] ) continue;
   R[u] = 1;
   BK( R, P & G[ u ], X & G[ u ] );
   R[u] = P[u] = 0, X[u] = 1;
public:
 void init( size_t n_ ) {
  n = n_{-};
  for ( size_t i = 0 ; i < n ; ++ i )
   G[ i ].reset();
  ans.reset();
 void add_edges( int u, bits S ) { G[ u ] = S; }
void add_edge( int u, int v ) {
  G[u][v] = G[v][u] = 1;
 int solve() {
  sort_by_degree(); // or simply iota( deo... )
  for ( size_t i = 0 ; i < n ; ++ i )
   deg[ i ] = G[ i ].count();
  bits pob, nob = 0; pob.set();
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
  for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t v = deo[ i ];</pre>
   bits tmp; tmp[ v ] = 1;
   BK( tmp, pob & G[ v ], nob & G[ v ] );
   pob[v] = 0, nob[v] = 1;
  return static_cast< int >( ans.count() );
};
3.10 MaxCliqueDyn
constexpr int kN = 150;
struct MaxClique { // Maximum Clique
 bitset<kN> a[kN], cs[kN];
int ans, sol[kN], q, cur[kN], d[kN], n;
 void init(int _n) {
  n = n, ans q = 0;
```

for (int i = 0; i < n; i++) a[i].reset();</pre>

fill(d[i+1], d[i+1]+n, inf);

```
void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
                                                                   for(int j=0; j<m; j++) {</pre>
                                                                    int v = e[j].v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
void csort(vector<int> &r, vector<int> &c)
 int mx = 1, km = max(ans - q + 1, 1), t = 0,
    m = int(r.size());
                                                                     d[i+1][u] = d[i][v]+e[j].c;
  cs[1].reset(); cs[2].reset();
                                                                     prv[i+1][u] = v;
  for (int i = 0; i < m; i++) {
                                                                     prve[i+1][u] = j;
  int p = r[i], k = 1;
   while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
                                                                  }
   cs[k][p] = 1;
  if (k < km) r[t++] = p;
                                                                 double solve(){
                                                                  // returns inf if no cycle, mmc otherwise
 c.resize(m);
                                                                  double mmc=inf;
  if (t) c[t - 1] = 0;
                                                                  int st = -1
  for (int k = km; k <= mx; k++) {</pre>
                                                                  bellman_ford();
  for (int p = int(cs[k]._Find_first());
                                                                  for(int i=0; i<n; i++) {</pre>
        < kN; p = int(cs[k]._Find_next(p))) {
                                                                   double avg=-inf;
                                                                   for(int k=0; k<n; k++) {</pre>
    r[t] = p; c[t++] = k;
                                                                    if(d[n][i]<inf-eps)</pre>
   }
 }
                                                                     avg=max(avg,(d[n][i]-d[k][i])/(n-k));
                                                                    else avg=max(avg,inf);
}
 void dfs(vector<int> &r, vector<int> &c, int 1,
                                                                   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
 bitset<kN> mask) {
 while (!r.empty()) {
   int p = r.back(); r.pop_back();
                                                                  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
  mask[p] = 0;
                                                                  for (int i=n; !vst[st]; st=prv[i--][st]) {
   if (q + c.back() <= ans) return;</pre>
                                                                   vst[st]++;
   cur[q++] = p;
                                                                   edgeID.PB(prve[i][st]);
   vector<int> nr, nc;
                                                                   rho.PB(st);
   bitset<kN> nmask = mask & a[p];
   for (int i : r)
                                                                  while (vst[st] != 2) {
    if (a[p][i]) nr.push_back(i);
                                                                   int v = rho.back(); rho.pop_back();
   if (!nr.empty()) {
                                                                   cycle.PB(v);
    if (1 < 4) {
                                                                   vst[v]++;
     for (int i : nr)
      d[i] = int((a[i] & nmask).count());
                                                                  reverse(ALL(edgeID));
     sort(nr.begin(), nr.end(),
                                                                  edgeID.resize(SZ(cycle));
      [&](int x, int y) {
  return d[x] > d[y];
                                                                  return mmc;
                                                                } mmc;
      });
  csort(nr, nc); dfs(nr, nc, l + 1, nmask);
} else if (q > ans) {
                                                                3.12 Minimum Steiner Tree
    ans = q; copy(cur, cur + q, sol);
                                                                // Minimum Steiner Tree
                                                                // 0(V 3^T + V^2 2^T)
   c.pop_back(); q--;
                                                                struct SteinerTree {
                                                                #define V 33
                                                                #define T 8
                                                                #define INF 1023456789
int solve(bitset<kN> mask) { // vertex mask
 vector<int> r, c;
                                                                 int n, dst[V][V], dp[1 << T][V], tdst[V];</pre>
 for (int i = 0; i < n; i++)
                                                                 void init(int _n) {
   if (mask[i]) r.push_back(i);
                                                                  n = _n;
 for (int i = 0; i < n; i++)
                                                                  for (int i = 0; i < n; i++) {</pre>
  d[i] = int((a[i] & mask).count());
                                                                   for (int j = 0; j < n; j++)
dst[i][j] = INF * (i != j);</pre>
 sort(r.begin(), r.end(),
  [&](int i, int j) { return d[i] > d[j]; });
  csort(r, c);
 dfs(r, c, 1, mask);
return ans; // sol[0 ~ ans-1]
                                                                 void add_edge(int ui, int vi, int wi) {
                                                                  dst[ui][vi] = min(dst[ui][vi], wi)
                                                                  dst[vi][ui] = min(dst[vi][ui], wi);
} graph;
                                                                 void shortest_path() {
3.11 Minimum Mean Cycle
                                                                  for (int k = 0; k < n; k++)
                                                                   for (int i = 0; i < n; i++)</pre>
/* minimum mean cycle O(VE) */
                                                                    for (int j = 0; j < n; j++)
struct MMC{
#define FZ(n) memset((n),0,sizeof(n))
                                                                     dst[i][j] = min(dst[i][j], dst[i][k] + dst[k][j]);
#define E 101010
#define V 1021
                                                                 int solve(const vector<int> &ter) {
#define inf 1e9
                                                                  int t = (int)ter.size();
                                                                  for (int i = 1; i < (1 << t); i++)
struct Edge { int v,u; double c; };
 int n, m, prv[V][V], prve[V][V], vst[V];
                                                                   fill_n(dp[i], n, INF);
Edge e[E];
                                                                  fill_n(dp[0], n, 0);
vector<int> edgeID, cycle, rho;
                                                                  for (int msk = 1; msk < (1 << t); msk++) {</pre>
double d[V][V];
                                                                   if (msk == (msk & (-msk))) {
void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
                                                                    int who = _{-}lg(msk);
                                                                    for (int i = 0; i < n; i++)
void add_edge( int vi , int ui , double ci )
{ e[ m ++ ] = { vi , ui , ci }; }
                                                                     dp[msk][i] = dst[ter[who]][i];
                                                                    continue:
void bellman_ford() {
 for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {</pre>
```

for (int i = 0; i < n; i++)

(submsk - 1) & msk)

for (int submsk = (msk - 1) & msk; submsk; submsk =

return cnt;

```
dp[msk][i] = min(dp[msk][i], dp[submsk][i] + dp[
    msk ^ submsk][i]);
                                                            };
   for (int i = 0; i < n; i++) {</pre>
                                                             4.2 Dijkstra Cost Flow
    tdst[i] = INF;
    for (int j = 0; j < n; j++)
                                                            // kN = #(vertices)
     tdst[i] = min(tdst[i], dp[msk][j] + dst[j][i]);
                                                             // MCMF.{Init, AddEdge, MincostMaxflow}
                                                             // MincostMaxflow(source, sink, flow_limit, &cost)
   copy_n(tdst, n, dp[msk]);
                                                                => flow
                                                             using Pii = pair<int, int>;
                                                             constexpr int kInf = 0x3f3f3f3f, kN = 500;
  int ans = INF;
 for (int i = 0; i < n; i++)
                                                             struct Edge {
  ans = min(ans, dp[(1 << t) - 1][i]);
                                                             int to, rev, cost, flow;
  return ans;
                                                            };
                                                             struct MCMF { // 0-based
} solver;
                                                             int n{}, m{}, s{}, t{};
                                                              vector<Edge> graph[kN];
      Mo's Algorithm on Tree
                                                              // Larger range for relabeling
                                                              int64_t dis[kN] = {}, h[kN] = {};
dfs u:
                                                              int p[kN] = {};
push u
                                                              void Init(int nn) {
 iterate subtree
                                                              n = nn;
push u
                                                               for (int i = 0; i < n; i++) graph[i].clear();</pre>
Let P = LCA(u, v) with St(u) <= St(v)
if (P == u) query[St(u), St(v)]
                                                              void AddEdge(int u, int v, int f, int c) {
else query[Ed(u), St(v)], query[St(P), St(P)]
                                                               graph[u].push_back({v,
                                                                static_cast<int>(graph[v].size()), c, f});
     Virtual Tree
3.14
                                                               graph[v].push_back(
vector<pair<int, int>> build(vector<int> vs, int r) {
                                                                {u, static_cast<int>(graph[u].size()) - 1,
vector<pair<int, int>> res;
                                                                 -c, 0});
sort(vs.begin(), vs.end(), [](int i, int j) {
 return dfn[i] < dfn[j]; });</pre>
                                                              bool Dijkstra(int &max_flow, int64_t &cost) {
vector < int > s = \{r\}
                                                               priority_queue<Pii, vector<Pii>, greater<>> pq;
for (int v : vs) if (v != r) {
                                                               fill_n(dis, n, kInf);
 if (int o = lca(v, s.back()); o != s.back()) {
                                                               dis[s] = 0:
   while (s.size() >= 2)
                                                               pq.emplace(0, s);
    if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
                                                               while (!pq.empty()) {
    res.emplace_back(s[s.size() - 2], s.back());
                                                                auto u = pq.top();
    s.pop_back();
                                                                pq.pop();
                                                                int v = u.second;
   if (s.back() != o) {
                                                                if (dis[v] < u.first) continue;</pre>
    res.emplace_back(s.back(), o);
                                                                for (auto &e : graph[v]) {
    s.back() = o;
                                                                 auto new_dis =
                                                                  dis[v] + e.cost + h[v] - h[e.to];
                                                                 if (e.flow > 0 && dis[e.to] > new_dis) {
  s.push_back(v);
                                                                  dis[e.to] = new_dis;
                                                                  p[e.to] = e.rev
for (size_t i = 1; i < s.size(); ++i)</pre>
                                                                  pq.emplace(dis[e.to], e.to);
 res.emplace_back(s[i - 1], s[i]);
 return res;
                                                               if (dis[t] == kInf) return false;
     Matching & Flow
                                                               for (int i = 0; i < n; i++) h[i] += dis[i];
                                                               int d = max_flow;
4.1 Bipartite Matching
                                                               for (int u = t; u != s;
struct BipartiteMatching {
                                                                  u = graph[u][p[u]].to) {
vector<int> X[N];
                                                                auto &e = graph[u][p[u]];
int fX[N], fY[N], n;
                                                                d = min(d, graph[e.to][e.rev].flow);
bitset<N> vis;
bool dfs(int x)
                                                              max_flow -= d;
  for (auto i : X[x]) if (not vis[i]) {
                                                               cost += int64_t(d) * h[t];
                                                              for (int u = t; u != s;
    u = graph[u][p[u]].to) {
  vis[i] = true;
if (fY[i] == -1 || dfs(fY[i])) {
   fY[fX[x] = i] = x;
                                                                auto &e = graph[u][p[u]];
                                                                e.flow += d;
    return true;
                                                               graph[e.to][e.rev].flow -= d;
 }
                                                              }
  return false;
                                                               return true;
 void init(int n_, int m) {
                                                              int MincostMaxflow(
                                                               int ss, int tt, int max_flow, int64_t &cost) {
 vis.reset();
 fill(X, X + (n = n_), vector<int>());
memset(fX, -1, sizeof(int) * n);
                                                               this->s = ss, this->t = tt;
                                                               cost = 0;
 memset(fY, -1, sizeof(int) * m);
                                                               fill_n(h, n, 0);
                                                               auto orig_max_flow = max_flow;
                                                               while (Dijkstra(max_flow, cost) && max_flow) {}
 void add_edge(int x, int y) { X[x].push_back(y); }
int solve() { // return how many pair matched
                                                               return orig_max_flow - max_flow;
 int cnt = 0;
  for (int i = 0; i < n; i++) {
                                                            };
  vis.reset();
                                                             4.3 Dinic
   cnt += dfs(i);
```

template <typename Cap = int64_t>

class Dinic{

```
private:
  struct E{
    int to, rev;
    Cap cap;
  };
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
       int u = bfs.front(); bfs.pop();
       for (auto e: G[u]) {
         if (e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
         bfs.push(e.to); lv[e.to] = lv[u] + 1;
    }
    return lv[ed] != -1;
  Cap DFS(int u, Cap f){
    if (u == ed) return f;
    Cap ret = 0:
     for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
       auto &e = G[u][i];
       if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
       Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
       if (f == 0) return ret;
    if (ret == 0) lv[u] = -1;
    return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
  G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
  st = st_, ed = ed_; Cap ret = 0;
    while (BFS()) {
       idx.assign(n, 0);
       Cap f = DFS(st, numeric_limits<Cap>::max());
       ret += f;
       if (f == 0) break;
     return ret;
  }
};
```

Flow Models

- · Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}$.
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect t o s with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f
 eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the
 - maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect t o s with capacity ∞ and let the flow from S to T be f'. If $f+f'\neq \sum_{v\in V, in(v)>0}in(v)$, there's no solution. Otherwise, f^{\prime} is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge \boldsymbol{e} on the graph.
- Construct minimum vertex cover from maximum matching ${\cal M}$ on bipartite graph(X,Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M, x \to y$ otherwise. 2. DFS from unmatched vertices in X.

 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect $x \to y$ with (cost,cap) = (c,1) if c>0, otherwise connect $y\to x$ with (cost, cap)=(-c,1)
 - 3. For each edge with $c\,<\,0$, sum these cost as K , then increase d(y)by 1, decrease $d(\boldsymbol{x})$ by 1
 - 4. For each vertex v with d(v)>0, connect S o v with (cost, cap)=(0, d(v))

- 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =
- 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum densitu induced subaraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \to v, v \in G$ with capacity K
 - 4. For each edge (u, v, w) in G, connect $u \to v$ and $v \to u$ with capacity
 - 5. For $v \in {\it G}$, connect it with sink $v \to t$ with capacity K + 2T $\left(\sum_{e \in E(v)} w(e)\right) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copu v', and connect $u' \to v'$ with weight w(u,v).
 - 2. Connect v
 ightarrow v' with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge
 - (v,t) with capacity $-p_v$. 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of
- 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

int fa[kN], pre[kN], match[kN], s[kN], v[kN];

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with capacity c_n
- 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4.5 General Graph Matching

namespace matching {

```
vector<int> g[kN];
queue<int> q;
void Init(int n) {
 for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
 for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
 g[u].push_back(v);
 g[v].push_back(u);
int Find(int u) {
return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
 static int tk = 0; tk++;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
void Blossom(int x, int y, int 1) {
  while (Find(x) != 1) {
  pre[x] = y, y = match[x];
  if (s[y] == 1) q.push(y), s[y] = 0;
  if (fa[x] == x) fa[x] = 1;
if (fa[y] == y) fa[y] = 1;
  x = pre[y];
 }
bool Bfs(int r, int n) {
 for (int i = 0; i \le n; ++i) fa[i] = i, s[i] = -1;
 while (!q.empty()) q.pop();
 q.push(r);
 s[r] = 0;
 while (!q.empty()) {
  int x = q.front(); q.pop();
  for (int u : g[x]) {
   if (s[u] == -1) {
pre[u] = x, s[u] = 1;
    if (match[u] == n) {
```

return rt;

```
4.8 Kuhn Munkres
     for (int a = u, b = x, last; b != n; a = last, b =
     pre[a])
                                                                 class KM {
      last = match[b], match[b] = a, match[a] = b;
                                                                 private:
     return true;
                                                                  static constexpr lld INF = 1LL << 60;</pre>
                                                                  vector<lld> hl,hr,slk;
    q.push(match[u]);
                                                                  vector<int> fl,fr,pre,qu;
    s[match[u]] = 0;
                                                                  vector<vector<lld>> w;
   } else if (!s[u] && Find(u) != Find(x)) {
                                                                  vector<bool> v1, vr;
    int 1 = LCA(u, x, n);
Blossom(x, u, 1);
                                                                  int n, q1, qr;
                                                                  bool check(int x) {
    Blossom(u, x, 1);
                                                                   if (v1[x] = true, f1[x] != -1)
                                                                    return vr[qu[qr++] = f1[x]] = true;
  }
                                                                   while (x != -1) swap(x, fr[fl[x] = pre[x]]);
                                                                   return false;
 return false;
                                                                  void bfs(int s) {
int Solve(int n) {
                                                                   fill(slk.begin(), slk.end(), INF);
 int res = 0;
                                                                   fill(v1.begin(), v1.end(), false);
fill(vr.begin(), vr.end(), false);
 for (int x = 0; x < n; ++x) {
 if (match[x] == n) res += Bfs(x, n);
                                                                   ql = qr = 0;
                                                                   vr[qu[qr++] = s] = true;
 return res;
                                                                   while (true) {
}}
                                                                    11d d;
                                                                    while (ql < qr) {</pre>
4.6 Global Min-Cut
                                                                     for (int x = 0, y = qu[ql++]; x < n; ++x) {
   if(!v1[x]&&s1k[x]>=(d=h1[x]+hr[y]-w[x][y])){
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
                                                                       if (pre[x] = y, d) slk[x] = d;
bool v[maxn], del[maxn];
                                                                        else if (!check(x)) return;
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
                                                                     }
                                                                     }
pair<int, int> phase(int n) {
                                                                    d = INF;
                                                                    for (int x = 0; x < n; ++x)
if (!v1[x] && d > s1k[x]) d = s1k[x];
memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
int s = -1, t = -1;
                                                                     for (int x = 0; x < n; ++x) {
 while (true) {
                                                                     if (v1[x]) h1[x] += d;
  int c = -1;
                                                                      else slk[x] -= d;
  for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
                                                                     if (vr[x]) hr[x] -= d;
   if (c == -1 \mid | g[i] > g[c]) c = i;
                                                                    for (int x = 0; x < n; ++x)
                                                                     if (!v1[x] && !slk[x] && !check(x)) return;
  if (c == -1) break;
 v[s = t, t = c] = true;
for (int i = 0; i < n; ++i) {
                                                                 public:
   if (del[i] || v[i]) continue;
                                                                  void init( int n_ ) {
   g[i] += w[c][i];
                                                                   qu.resize(n = n_);
                                                                   fl.assign(n, -1); fr.assign(n, -1);
hr.assign(n, 0); hl.resize(n);
  }
                                                                   w.assign(n, vector<lld>(n));
 return make_pair(s, t);
                                                                   slk.resize(n); pre.resize(n);
                                                                   vl.resize(n); vr.resize(n);
int mincut(int n) {
int cut = 1e9:
 memset(del, false, sizeof(del));
                                                                  void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 for (int i = 0; i < n - 1; ++i) {
                                                                  11d solve() {
                                                                   for (int i = 0; i < n; ++i)
 int s, t; tie(s, t) = phase(n);
  del[t] = true; cut = min(cut, g[t]);
                                                                    hl[i] = *max_element(w[i].begin(), w[i].end());
  for (int j = 0; j < n; ++j) {
                                                                   for (int i = 0; i < n; ++i) bfs(i);</pre>
                                                                   11d res = 0;
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
  }
                                                                   for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
                                                                   return res;
                                                                  }
 return cut;
                                                                } km;
4.7 GomoryHu Tree
                                                                 4.9 Minimum Cost Circulation
int g[maxn];
                                                                 struct Edge { int to, cap, rev, cost; };
                                                                 vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
bool mark[kN];
vector<edge> GomoryHu(int n){
 vector<edge> rt;
 for(int i=1;i<=n;++i)g[i]=1;</pre>
 for(int i=2;i<=n;++i){</pre>
                                                                 int NegativeCycle(int n) {
                                                                  memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
  int t=g[i];
  flow.reset(); // clear flows on all edge
  rt.push_back({i,t,flow(i,t)});
                                                                  int upd = -1;
  flow.walk(i); // bfs points that connected to i (use
                                                                  for (int i = 0; i <= n; ++i) {
    edges not fully flow)
                                                                   for (int j = 0; j < n; ++j) {
  for(int j=i+1;j<=n;++j){</pre>
                                                                    int idx = 0;
                                                                    for (auto &e : g[j]) {
   if(g[j]==t && flow.connect(j))g[j]=i; // check if i
    can reach i
                                                                     if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
  }
                                                                       dist[e.to] = dist[j] + e.cost;
```

pv[e.to] = j, ed[e.to] = idx;

if (i == n) {
 upd = j;

```
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      while(!mark[upd])mark[upd]=1,upd=pv[upd];
     return upd:
     }
    idx++:
 }
return -1;
int Solve(int n) {
int rt = -1, ans = 0;
while ((rt = NegativeCycle(n)) >= 0) {
 memset(mark, false, sizeof(mark));
 vector<pair<int, int>> cyc;
 while (!mark[rt]) {
  cyc.emplace_back(pv[rt], ed[rt]);
  mark[rt] = true;
  rt = pv[rt];
 reverse(cyc.begin(), cyc.end());
 int cap = kInf;
 for (auto &i : cyc)
  auto &e = g[i.first][i.second];
  cap = min(cap, e.cap);
 for (auto &i : cyc)
  auto &e = g[i.first][i.second];
  e.cap -= cap;
  g[e.to][e.rev].cap += cap;
  ans += e.cost * cap;
 }
return ans;
4.10 Minimum Cost Maximum Flow
class MiniCostMaxiFlow{
using Cap = int; using Wei = int64_t;
using PCW = pair<Cap,Wei>;
static constexpr Cap INF_CAP = 1 << 30;</pre>
static constexpr Wei INF_WEI = 1LL<<60;</pre>
private:
struct Edge{
```

```
int to, back;
Cap cap; Wei wei;
Edge() {}
Edge(int a,int b, Cap c, Wei d):
  to(a),back(b),cap(c),wei(d) {}
};
int ori, edd;
vector<vector<Edge>> G;
vector<int> fa, wh;
vector<bool> inq;
vector<Wei> dis;
PCW SPFA(){
 fill(inq.begin(),inq.end(),false);
 fill(dis.begin(), dis.end(), INF_WEI);
 queue<int> qq; qq.push(ori);
 dis[ori] = 0;
 while(not qq.empty()){
  int u=qq.front();qq.pop();
  ing[u] = false;
  for(int i=0;i<SZ(G[u]);++i){</pre>
   Edge e=G[u][i];
   int v=e.to; Wei d=e.wei;
   if(e.cap <= 0 | |dis[v] <= dis[u] + d)
    continue
   dis[v] = dis[u] + d;
   fa[v] = u, wh[v] = i;
   if (ing[v]) continue;
   qq.push(v);
   inq[v] = true;
  }
 if(dis[edd]==INF_WEI) return {-1, -1};
 Cap mw=INF_CAP;
 for(int i=edd;i!=ori;i=fa[i])
 mw=min(mw,G[fa[i]][wh[i]].cap);
 for (int i=edd;i!=ori;i=fa[i]){
  auto &eg=G[fa[i]][wh[i]];
  eg.cap -= mw;
```

```
G[eg.to][eg.back].cap+=mw;
  return {mw, dis[edd]};
public:
 void init(int n){
  G.clear();G.resize(n);
  fa.resize(n);wh.resize(n);
  inq.resize(n); dis.resize(n);
 void add_edge(int st, int ed, Cap c, Wei w){
  G[st].emplace_back(ed,SZ(G[ed]),c,w);
  G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
 PCW solve(int a, int b){
 ori = a, edd = b;
Cap cc=0; Wei ww=0;
  while(true){
  PCW ret=SPFA();
   if(ret.first==-1) break;
   cc+=ret.first:
   ww+=ret.first * ret.second;
  return {cc,ww};
} mcmf;
4.11 Maximum Weight Graph Matching
struct WeightGraph {
 static const int inf = INT_MAX;
 static const int maxn = 514;
 struct edge {
  int u, v, w;
  edge(){}
  edge(int u, int v, int w): u(u), v(v), w(w) {}
 int n, n_x;
 edge g[maxn * 2][maxn * 2];
 int lab[maxn * 2];
 int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa
    [maxn * 2];
 int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
    maxn * 2]
 vector<int> flo[maxn * 2];
 queue<int> q;
 int e_delta(const edge &e) { return lab[e.u] + lab[e.v
    ] - g[e.u][e.v].w * 2; }
 void update_slack(int u, int x) { if (!slack[x] ||
    e_delta(g[u][x]) < e_delta(g[slack[x]][x])) slack[x
    ] = u; }
 void set_slack(int x) {
  slack[x] = 0;
  for (int u = 1; u <= n; ++u)</pre>
   if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
    update_slack(u, x);
 void q_push(int x) {
  if (x \le n) q.push(x);
  else for (size_t i = 0; i < flo[x].size(); i++)</pre>
    q_push(flo[x][i]);
 void set_st(int x, int b) {
  st[x] = b;
  if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)
     set_st(flo[x][i], b);
 int get_pr(int b, int xr) {
  int pr = find(flo[b].begin(), flo[b].end(), xr) - flo
    [b].begin();
  if (pr % 2 == 1)
   reverse(flo[b].begin() + 1, flo[b].end());
   return (int)flo[b].size() - pr;
  return pr;
 void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u, xr)
  for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo</pre>
    [u][i ^ 1]);
```

```
set_match(xr, v);
 rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
                                                               }
   end());
void augment(int u, int v) {
 for (; ; ) {
  int xnv = st[match[u]];
  set_match(u, v);
  if (!xnv) return;
  set_match(xnv, st[pa[xnv]]);
  u = st[pa[xnv]], v = xnv;
}
int get_lca(int u, int v) {
 static int t = 0;
 for (++t; u || v; swap(u, v)) {
  if (u == 0) continue;
  if (vis[u] == t) return u;
  vis[u] = t;
  u = st[match[u]];
  if (u) u = st[pa[u]];
 }
 return 0;
void add_blossom(int u, int lca, int v) {
 int b = n + 1;
 while (b <= n_x && st[b]) ++b;</pre>
 if (b > n_x) ++n_x;
 lab[b] = 0, S[b] = 0
 match[b] = match[lca];
 flo[b].clear();
 flo[b].push_back(lca);
 for (int x = u, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
 reverse(flo[b].begin() + 1, flo[b].end())
 for (int x = v, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
 set_st(b, b);
 for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].w
 for (int x = 1; x \le n; ++x) flo_from[b][x] = 0;
 for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  for (int x = 1; x <= n_x; ++x)
   if (g[b][x].w == 0 \mid \mid e_delta(g[xs][x]) < e_delta(g[xs][x])
   [b][x]))
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
  for (int x = 1; x <= n; ++x)
   if (flo_from[xs][x]) flo_from[b][x] = xs;
 set_slack(b);
void expand_blossom(int b) {
 for (size_t i = 0; i < flo[b].size(); ++i)</pre>
  set_st(flo[b][i], flo[b][i]);
 int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
   xr);
 for (int i = 0; i < pr; i += 2)
  int xs = flo[b][i], xns = flo[b][i + 1];
  pa[xs] = g[xns][xs].u;
  S[xs] = 1, S[xns] = 0;
  slack[xs] = 0, set_slack(xns);
  q_push(xns);
 S[xr] = 1, pa[xr] = pa[b];
 for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  S[xs] = -1, set_slack(xs);
 st[b] = 0;
bool on_found_edge(const edge &e) {
 int u = st[e.u], v = st[e.v];
 if (S[v] == -1) {
  pa[v] = e.u, S[v] = 1
  int nu = st[match[v]]
  slack[v] = slack[nu] = 0;
 S[nu] = 0, q_push(nu);
} else if (S[v] == 0) {
  int lca = get_lca(u, v);
```

```
if (!lca) return augment(u,v), augment(v,u), true;
  else add_blossom(u, lca, v);
 return false;
bool matching() {
 \begin{array}{lll} memset(S+1, -1, \ sizeof(int) \ * \ n_x); \\ memset(slack + 1, \ \theta, \ sizeof(int) \ * \ n_x); \end{array} 
 q = queue<int>();
for (int x = 1; x <= n_x; ++x)
if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0,
   q_push(x);
 if (q.empty()) return false;
  while (q.size()) {
   int u = q.front(); q.pop();
   if (S[st[u]] == 1) continue
   for (int v = 1; v <= n; ++v)
    if (g[u][v].w > 0 && st[u] != st[v]) {
     if (e_delta(g[u][v]) == 0) {
      if (on_found_edge(g[u][v])) return true;
      } else update_slack(u, st[v]);
  int d = inf;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b && S[b] == 1) d = min(d, lab[b] / 2)
  for (int x = 1; x <= n_x; ++x)
if (st[x] == x && slack[x]) {</pre>
    if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x])
    else if (S[x] == 0) d = min(d, e_delta(g[slack[x
   ]][x]) / 2);
  for (int u = 1; u <= n; ++u) {
   if (S[st[u]] == 0) {
    if (lab[u] <= d) return 0;</pre>
    lab[u] -= d;
   } else if (S[st[u]] == 1) lab[u] += d;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b) {
    if (S[st[b]] == 0) lab[b] += d * 2;
    else if (S[st[b]] == 1) lab[b] -= d * 2;
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
if (st[x] == x && slack[x] && st[slack[x]] != x &&</pre>
   e_{delta}(g[slack[x]][x]) == 0)
    if (on_found_edge(g[slack[x]][x])) return true;
  for (int b = n + 1; b <= n_x; ++b)
if (st[b] == b && S[b] == 1 && lab[b] == 0)
   expand_blossom(b);
 return false;
pair<long long, int> solve() {
memset(match + 1, 0, sizeof(int) * n);
 n_x = n:
 int n_matches = 0;
 long long tot_weight = 0;
 for (int u = 0; u \le n; ++u) st[u] = u, flo[u].clear
   ();
 int w_max = 0;
 for (int u = 1; u <= n; ++u)</pre>
 for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);</pre>
   w_max = max(w_max, g[u][v].w);
 for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
 while (matching()) ++n_matches;
for (int u = 1; u <= n; ++u)
  if (match[u] && match[u] < u)</pre>
   tot_weight += g[u][match[u]].w;
 return make_pair(tot_weight, n_matches);
void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
   g[vi][ui].w = wi; }
void init(int _n) {
 n = _n;
for (int u = 1; u <= n; ++u)
```

```
for (int v = 1; v <= n; ++v)
  g[u][v] = edge(u, v, 0);
}
};</pre>
```

5 Math

5.1 Strling Number

5.1.1 First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1) \dots (x+n-1) = \sum_{k=0}^n S_1(n,k) x^k$$

$$g(x) = x(x+1) \dots (x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.1.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

$$S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}$$

5.2 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else exgcd(y,x%y,g,b,a),b==(x/y)*a;
}
```

5.3 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
vector<T> d(output.size() + 1), me, he;
for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
 for (size_t j = 0; j < me.size(); ++j)</pre>
  d[i] += output[i - j - 2] * me[j];
  if ((d[i] -= output[i - 1]) == 0) continue;
 if (me.empty()) {
  me.resize(f = i);
  continue;
 vector<T> o(i - f - 1);
 T k = -d[i] / d[f]; o.push_back(-k);
 for (T x : he) o.push_back(x * k);
 if (o.size() < me.size()) o.resize(me.size());</pre>
 for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];
 if (i-f+he.size() >= me.size()) he = me, f = i;
 me = o:
return me;
```

5.4 Charateristic Polynomial

```
vector<vector<int>>> &A) {
  int N = A.size();
  vector<vector<int>>> H = A;
  for (int i = 0; i < N - 2; ++i) {
    if (!H[i + 1][i]) {
      for (int j = i + 2; j < N; ++j) {
        if (H[j][i]) {
          for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j][k]);
          for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k][j][i]);
          break;
      }
    }
}</pre>
```

```
if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
   * H[i + 1][k] * (kP - coef)) % kP;
for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
 return H;
vector<int> CharacteristicPoly(const vector<vector<int</pre>
    >> &A) {
 int N = A.size();
 auto H = Hessenberg(A);
 for (int i = 0; i < N; ++i) {
  for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
 vector<vector<int>> P(N + 1, vector<int>(N + 1));
 P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
  P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
    1];
  int val = 1;
  for (int j = i - 1; j >= 0; --j) {
  int coef = 1LL * val * H[j][i - 1] % kP;
   for (int k = 0; k <= j; ++k) P[i][k] = (P[i][k] + 1
LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 if (N & 1) {
  for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
 return P[N];
```

5.5 Chinese Remainder

x = a1 % m1

x = a2 % m2

```
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
// 0 <= x < lcm(m1, m2)
5.6 De-Bruijn
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
  if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
  else {
  aux[t] = aux[t - p];
  db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {
  aux[t] = i:
   db(t + 1, t, n, k);
int de_bruijn(int k, int n) {
 // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
 if (k == 1) {
  res[0] = 0;
  return 1;
```

for (int i = 0; i < k * n; i++) aux[i] = 0;

5.7 DiscreteLog

db(1, 1, n, k);

return sz;

```
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
```

```
// x^? \equiv y (mod M)
Int t = 1, c = 0, g = 1;
for (Int M_ = M; M_ > 0; M_ >>= 1)
    g = g * x % M;
for (g = gcd(g, M); t % g != 0; ++c) {
    if (t == y) return c;
    t = t * x % M;
}
if (y % g != 0) return -1;
t /= g, y /= g, M /= g;
Int h = 0, gs = 1;
for (; h * h < M; ++h) gs = gs * x % M;
unordered_map<Int, Int> bs;
for (Int s = 0; s < h; bs[y] = ++s)
    y = y * x % M;
for (Int s = 0; s < M; s += h) {
    t = t * gs % M;
    if (bs.count(t)) return c + s + h - bs[t];
}
return -1;
}</pre>
```

5.8 Extended Euler

$$a^b \equiv \begin{cases} a^{(b \mod \varphi(m)) + \varphi(m)} & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod{m}$$

5.9 ExtendedFloorSum

```
g(a,b,c,n) = \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor
                              \left( \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \right)
                               +g(a \bmod c, b \bmod c, c, n),
                                                                                                                      a \geq c \vee b \geq c
                                                                                                                      n < 0 \lor a = 0
                               \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
                              -h(c, c-b-1, a, m-1)),
                                                                                                                       otherwise
h(a,b,c,n) = \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2
                              \left( \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right)
                                +\lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1)
                                +h(a \bmod c, b \bmod c, c, n)
                                +2\lfloor \tfrac{a}{c}\rfloor \cdot g(a \bmod c, b \bmod c, c, n)
                                +2\lfloor \frac{\bar{b}}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
                                                                                                                      a \geq c \vee b \geq c
                                                                                                                       n<0\vee a=0
```

nm(m+1) - 2g(c, c-b-1, a, m-1)

-2f(c, c-b-1, a, m-1) - f(a, b, c, n), otherwise

5.10 Fast Fourier Transform

```
const int mod = 1000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
  static_assert (M1 <= M2 && M2 <= M3);</pre>
  constexpr int64_t r12 = modpow(M1, M2-2, M2);
 constexpr int64_t r13 = modpow(M1, M3-2, M3);
  constexpr int64_t r23 = modpow(M2, M3-2, M3);
  constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
 B = (B - A + M2) * r12 % M2;

C = (C - A + M3) * r13 % M3;
 C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
for (int i = 0; i <= maxn; i++)</pre>
  omega[i] = cplx(cos(2 * pi * j / maxn),
     sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
  int z = __builtin_ctz(n) - 1;
for (int i = 0; i < n; ++i) {
  int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^{=(i >> j & 1)<<(z - j);
  if (x > i) swap(v[x], v[i]);
```

```
for (int s = 2; s <= n; s <<= 1) {
  int z = s \gg 1;
  for (int i = 0; i < n; i += s) {
   for (int k = 0; k < z; ++k) {
   cplx x = v[i + z + k] * omega[maxn / s * k];
     v[i + z + k] = v[i + k] - x;
     v[i+k] = v[i+k] + x;
void ifft(vector<cplx> &v, int n) {
  fft(v, n); reverse(v.begin() + 1, v.end());
  for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);</pre>
VL convolution(const VI &a, const VI &b) {
 // Should be able to handle N <= 10^5, C <= 10^4
 int sz = 1;
 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 vector<cplx> v(sz);
 for (int i = 0; i < sz; ++i) {
  double re = i < a.size() ? a[i] : 0;</pre>
  double im = i < b.size() ? b[i] : 0;</pre>
  v[i] = cplx(re, im);
 fft(v, sz);
 for (int i = 0; i <= sz / 2; ++i) {
  int j = (sz - i) & (sz - 1);
cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
     * cplx(0, -0.25);
  if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
     ].conj()) * cplx(0, -0.25);
   v[i] = x;
 ifft(v, sz);
 VL c(sz);
 for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
VI convolution_mod(const VI &a, const VI &b, int p) {
 int sz = 1;
 while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
 vector<cplx> fa(sz), fb(sz);
for (int i = 0; i < (int)a.size(); ++i)
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 for (int i = 0; i < (int)b.size(); ++i)</pre>
  fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
 fft(fa, sz), fft(fb, sz);
 double r = 0.25 / sz;
 cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
  int j = (sz - i) & (sz - 1);
  cplx a1 = (fa[i] + fa[j].conj());
cplx a2 = (fa[i] - fa[j].conj()) * r2;
  cplx b1 = (fb[i] + fb[j].conj()) * r3;
cplx b2 = (fb[i] - fb[j].conj()) * r4;
  if (i != j) {
   cplx c1 = (fa[j] + fa[i].conj());
    cplx c2 = (fa[j] - fa[i].conj()) * r2;
    cplx d1 = (fb[j] + fb[i].conj()) * r3;
cplx d2 = (fb[j] - fb[i].conj()) * r4;
    fa[i] = c1 * d1 + c2 * d2 * r5;
    fb[i] = c1 * d2 + c2 * d1;
  fa[j] = a1 * b1 + a2 * b2 * r5;
  fb[j] = a1 * b2 + a2 * b1;
 fft(fa, sz), fft(fb, sz);
 vector<int> res(sz);
 for (int i = 0; i < sz; ++i) {
  long long a = round(fa[i].re), b = round(fb[i].re),
         c = round(fa[i].im);
  res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
 }
 return res;
5.11 FloorSum
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
```

// @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64

```
1lu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
                                                                 if (witn(mpow(m % x, odd, x), odd, x, t))
 llu ans = 0;
                                                                   return false:
 while (true)
                                                                return true;
  if (a >= m) {
   ans += n * (n - 1) / 2 * (a / m); a %= m;
                                                               5.14 NTT
  if (b >= m) {
                                                               template <int mod, int G, int maxn>
   ans += n * (b / m); b %= m;
                                                               struct NTT {
                                                                static_assert (maxn == (maxn & -maxn));
  llu y_max = a * n + b;
                                                                int roots[maxn];
 if (y_max < m) break;</pre>
                                                                NTT () {
  // y_max < m * (n + 1)
                                                                 int r = modpow(G, (mod - 1) / maxn);
  // floor(y_max / m) <= n
                                                                 for (int i = maxn >> 1; i; i >>= 1) {
  n = (11u)(y_max / m), b = (11u)(y_max % m);
                                                                   roots[i] = 1;
  swap(m, a);
                                                                   for (int j = 1; j < i; j++)
roots[i + j] = modmul(roots[i + j - 1], r);</pre>
 return ans;
                                                                   r = modmul(r, r);
11d floor_sum(lld n, lld m, lld a, lld b) {
 llu ans = 0;
                                                                // n must be 2^k, and 0 \le F[i] < mod
 if (a < 0) {
                                                                void operator()(int F[], int n, bool inv = false) {
 11u a2 = (a \% m + m) \% m;
                                                                 for (int i = 0, j = 0; i < n; i++) {
 ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
                                                                   if (i < j) swap(F[i], F[j]);</pre>
  a = a2:
                                                                   for (int k = n > 1; (j^k < k; k > = 1);
 if (b < 0) {
                                                                 for (int s = 1; s < n; s *= 2) {
 11\dot{u} b2 = (b % m + m) % m;
                                                                   for (int i = 0; i < n; i += s * 2) {
  ans -= 1ULL * n * ((b2 - b) / m);
                                                                   for (int j = 0; j < s; j++) {
 b = b2:
                                                                     int a = F[i+j]
                                                                     int b = modmul(F[i+j+s], roots[s+j]);
return ans + floor_sum_unsigned(n, m, a, b);
                                                                    F[i+j] = modadd(a, b); // a + b

F[i+j+s] = modsub(a, b); // a - b
5.12 FWT
                                                                   }
/* xor convolution:
* x = (x0, x1) , y = (y0, y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                 if (inv) {
                                                                   int invn = modinv(n);
                                                                   for (int i = 0; i < n; i++)
 * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
                                                                   F[i] = modmul(F[i], invn);
                                                                   reverse(F + 1, F + n);
 *z = (1/2) *z'
 * or convolution:
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
 * and convolution:
                                                               NTT<2013265921, 31, 1048576> ntt;
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
const LL MOD = 1e9+7;
                                                               5.15 Partition Number
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
                                                               int b = sqrt(n);
 for( int d = 1 ; d < N ; d <<= 1 ) {</pre>
                                                               ans[0] = tmp[0] = 1;
  int d2 = d << 1;
                                                               for (int i = 1; i <= b; i++) {
  for( int s = 0; s < N; s += d2)
                                                                for (int rep = 0; rep < 2; rep++)</pre>
   for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
LL ta = x[ i ] , tb = x[ j ];</pre>
                                                                 for (int j = i; j <= n - i * i; j++)
                                                                modadd(tmp[j], tmp[j-i]);
for (int j = i * i; j <= n; j++)</pre>
    x[i] = ta+tb:
    x[ j ] = ta-tb;
if( x[ i ] >= MOD ) x[ i ] -= MOD;
                                                                 modadd(ans[j], tmp[j - i * i]);
    if (x[j] < 0) x[j] += MOD;
                                                               5.16 Pi Count (Linear Sieve)
 if( inv )
                                                               static constexpr int N = 1000000 + 5;
  for( int i = 0 ; i < N ; i++ ) {
                                                               11d pi[N];
   x[ i ] *= inv( N, MOD );
                                                               vector<int>
                                                                            primes;
   x[ i ] %= MOD;
                                                               bool sieved[N];
                                                               1ld cube_root(lld x){
}
                                                                1ld s=cbrt(x-static_cast<long double>(0.1));
                                                                while(s*s*s <= x) ++s;</pre>
5.13 Miller Rabin
                                                                return s-1;
bool isprime(llu x) {
  static auto witn = [](llu a, llu u, llu n, int t) {
                                                               11d square_root(11d x){
  if (!a) return false;
                                                                1ld s=sqrt(x-static_cast<long double>(0.1));
  while (t--)
                                                                while(s*s <= x) ++s;
   llu a2 = mmul(a, a, n);
                                                                return s-1;
   if (a2 == 1 && a != 1 && a != n - 1) return true;
   a = a2;
                                                               void init(){
  }
                                                                primes.reserve(N);
  return a != 1;
                                                                primes.push_back(1);
                                                                for(int i=2;i<N;i++) {</pre>
                                                                 if(!sieved[i]) primes.push_back(i);
 if (x < 2) return false;</pre>
 if (!(x & 1)) return x == 2;
                                                                 pi[i] = !sieved[i] + pi[i-1];
 int t = __builtin_ctzll(x - 1);
                                                                 for(int p: primes) if(p > 1) {
                                                                   if(p * i >= N) break;
 llu odd = (x - 1) >> t;
 for (llu m:
                                                                   sieved[p * i] = true;
  {2, 325, 9375, 28178, 450775, 9780504, 1795265022})
                                                                   if(p % i == 0) break;
```

```
Poly X = Poly(*this, (size() + 1) / 2).Inv().isz(sz),
                                                               Y(*this, sz);
ntt(X.data(), sz), ntt(Y.data(), sz);
}
                                                               fi(0, sz) X[i] = modmul(X[i], modsub(2, modmul(X[i],
11d phi(11d m, 11d n) {
static constexpr int MM = 80000, NN = 500;
                                                                  Y[i])));
static lld val[MM][NN];
                                                               ntt(X.data(), sz, true);
                                                               return X.isz(size());
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
if(n == 0) return m;
if(primes[n] >= m) return 1;
                                                              Poly Sqrt() const { // coef[0] \in [1, mod)^2
                                                               if (size() == 1) return V{QuadraticResidue((*this)
lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
if(m < MM\&n < NN) val[m][n] = ret+1;
                                                                  [0], mod)};
                                                               Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
return ret;
                                                                  size()):
11d pi_count(11d);
                                                                return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
11d P2(11d m, 11d n) {
                                                                  + 1);
1ld sm = square_root(m), ret = 0;
for(lld i = n+1;primes[i]<=sm;i++)</pre>
                                                              pair<Poly, Poly> DivMod(const Poly &rhs) const {
                                                               if (size() < rhs.size()) return {V{0}, *this};</pre>
 ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
                                                               const int sz = size() - rhs.size() + 1;
return ret;
                                                               Poly X(rhs); X.irev().isz(sz);
11d pi_count(11d m) {
                                                               Poly Y(*this); Y.irev().isz(sz);
                                                               Poly Q = Y.Mul(X.Inv()).isz(sz).irev();
if(m < N) return pi[m];</pre>
                                                               X = rhs.Mul(Q), Y = *this;
fi(0, size()) Y[i] = modsub(Y[i], X[i]);
11d n = pi_count(cube_root(m));
return phi(m, n) + n - 1 - P2(m, n);
                                                               return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
5.17 Pollard Rho
                                                              Poly Dx() const {
// does not work when n is prime
                                                               Poly ret(size() - 1);
// return any non-trivial factor
                                                               fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
1lu pollard_rho(llu n) {
                                                                 1]);
static auto f = [](llu x, llu k, llu m) {
                                                               return ret.isz(max<int>(1, ret.size()));
    return add(k, mul(x, x, m), m); };
if (!(n & 1)) return 2;
                                                              Poly Sx() const {
mt19937 rnd(120821011);
                                                               Poly ret(size() + 1);
                                                               fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
while (true) {
 11u y = 2, yy = y, x = rnd() % n, t = 1;
                                                                  this)[i]);
 for (llu sz = 2; t == 1; sz <<= 1, y = yy) {
                                                               return ret;
  for (llu i = 0; t == 1 && i < sz; ++i) {
   yy = f(yy, x, n);
                                                              Poly Ln() const { // coef[0] == 1
                                                               return Dx().Mul(Inv()).Sx().isz(size());
    t = gcd(yy > y ? yy - y : y - yy, n);
                                                              Poly Exp() const { // coef[0] == 0
  if (size() == 1) return V{1};
  if (t != 1 && t != n) return t;
                                                               Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size)
                                                               ());
Poly Y = X.Ln(); Y[0] = mod - 1;
     Polynomial Operations
5.18
                                                               fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
                                                               return X.Mul(Y).isz(size());
using V = vector<int>;
#define fi(1, r) for (int i = int(1); i < int(r); ++i)
template <int mod, int G, int maxn> struct Poly : V {
                                                              Poly Pow(const string &K) const {
                                                               int nz = 0;
static uint32_t n2k(uint32_t n) {
 if (n <= 1) return 1;
                                                               while (nz < size() && !(*this)[nz]) ++nz;</pre>
  return 1u << (32 - __builtin_clz(n - 1));</pre>
                                                               int nk = 0, nk2 = 0;
                                                               for (char c : K) {
                                                                nk = (nk * 10 + c - '0') % mod;
static NTT<mod,G,maxn> ntt; // coefficients in [0, P)
                                                                nk2 = nk2 * 10 + c - '0';
explicit Poly(int n = 1) : V(n) {}
Poly(const V &v) : V(v) {}
                                                                 if (nk2 * nz >= size())
                                                                 return Poly(size());
Poly(const Poly &p, size_t n) : V(n) {
                                                                nk2 \% = mod - 1;
 copy_n(p.data(), min(p.size(), n), data());
                                                               if (!nk && !nk2) return Poly(V{1}, size());
Poly &irev() { return reverse(data(), data() + size())
                                                               Poly X = V(data() + nz, data() + size() - nz * (nk2 -
    , *this; }
                                                                   1));
Poly &isz(int sz) { return resize(sz), *this; }
Poly &iadd(const Poly &rhs) { // n() == rhs.n()
                                                               int x0 = X[0];
                                                                return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
 fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
  return *this;
                                                                  modpow(x0, nk2)).irev().isz(size()).irev();
                                                              Poly InvMod(int L) { // (to evaluate linear recursion)
Poly &imul(int k) {
                                                               Poly R{1, 0}; // *this * R mod x^L = 1 (*this[0] ==
 fi(0, size())(*this)[i] = modmul((*this)[i], k);
  return *this;
                                                                 1)
                                                               for (int level = 0; (1 << level) < L; ++level)</pre>
Poly Mul(const Poly &rhs) const {
                                                                Poly 0 = R.Mul(Poly(data(), min<int>(2 << level,
                                                                  size())));
 const int sz = n2k(size() + rhs.size() - 1);
 Poly X(*this, sz), Y(rhs, sz);
ntt(X.data(), sz), ntt(Y.data(), sz);
fi(0, sz) X[i] = modmul(X[i], Y[i]);
                                                                 Poly Q(2 << level); Q[0] = 1;
                                                                for (int j = (1 << level); j < (2 << level); ++j)</pre>
                                                                 Q[j] = modsub(mod, O[j]);
 ntt(X.data(), sz, true);
return X.isz(size() + rhs.size() - 1);
                                                                R = R.Mul(Q).isz(4 << level);
                                                               return R.isz(L);
Poly Inv() const { // coef[0] != 0
 if (size() == 1) return V{modinv(*begin())};
                                                              static int LinearRecursion(const V &a, const V &c,
 const int sz = n2k(size() * 2);
```

 $int64_t n) { // a_n = \sum c_j a_(n-j)}$

if (d[i][s] < eps) continue;
if (r == -1 || \</pre>

d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;

```
const int k = (int)a.size();
                                                                    if (r == -1) return false;
  assert((int)c.size() == k + 1);
  Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
                                                                    pivot(r, s);
  fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
  C[k] = 1;
                                                                  VD solve(const VVD &a, const VD &b, const VD &c) {
  while (n)
   if (n % 2) W = W.Mul(M).DivMod(C).second;
                                                                   m = b.size(), n = c.size();
   n /= 2, M = M.Mul(M).DivMod(C).second;
                                                                   d = VVD(m + 2, VD(n + 2))
                                                                   for (int i = 0; i < m; ++i)</pre>
                                                                    for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
  int ret = 0;
  fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
                                                                   p.resize(m), q.resize(n + 1);
                                                                   for (int i = 0; i < m; ++i)
  return ret:
                                                                    p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
                                                                   for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
                                                                   q[n] = -1, d[m + 1][n] = 1;
#undef fi
using Poly_t = Poly<998244353, 3, 1 << 20>;
                                                                   int r = 0;
                                                                   for (int i = 1; i < m; ++i)
if (d[i][n + 1] < d[r][n + 1]) r = i;
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
5.19 Quadratic residue
                                                                   if (d[r][n + 1] < -eps) {</pre>
struct S {
                                                                    pivot(r, n);
 int MOD, w;
                                                                    if (!phase(1) || d[m + 1][n + 1] < -eps)</pre>
 int64_t x, y;
                                                                     return VD(n, -inf);
                                                                    for (int i = 0; i < m; ++i) if (p[i] == -1) {
 S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
  : MOD(m), w(w_{-}), x(x_{-}), y(y_{-}) {}
                                                                     int s = min_element(d[i].begin(), d[i].end() - 1)
 S operator*(const S &rhs) const {
                                                                          - d[i].begin();
  int w_ = w;
if (w_ == -1) w_ = rhs.w;
                                                                     pivot(i, s);
                                                                    }
  assert(w_! = -1 \text{ and } w_ == rhs.w);
  return { MOD, w_,
(x * rhs.x + y * rhs.y % MOD * w) % MOD,
                                                                   if (!phase(0)) return VD(n, inf);
                                                                   VD x(n);
   (x * rhs.y + y * rhs.x) % MOD };
                                                                   for (int i = 0; i < m; ++i)</pre>
 }
                                                                   if (p[i] < n) x[p[i]] = d[i][n + 1];
                                                                   return x;
                                                                  }}
int get_root(int n, int P) {
 if (P == 2 or n == 0) return n;
                                                                  5.21 Simplex Construction
  if (qpow(n, (P - 1) / 2, P) != 1) return -1;
  auto check = [&](int x) {
                                                                  Standard form: maximize \sum_{1 \le i \le n} c_i x_i such that for all 1 \le j \le m,
                                                                  \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j and x_i \geq 0 for all 1 \leq i \leq n.
    return qpow(x, (P - 1) / 2, P); };
  if (check(n) == P-1) return -1
  int64_t a; int w; mt19937 rnd(7122);
                                                                    1. In case of minimization, let c'_i = -c_i
  do { a = rnd() % P;
    w = ((a * a - n) \% P + P) \% P;
                                                                    2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
  } while (check(w) != P - 1);
  return qpow(S(P, w, a, 1), (P + 1) / 2).x;
                                                                    3. \sum_{1 \le i \le n} A_{ji} x_i = b_j
                                                                          • \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
5.20 Simplex
                                                                          • \sum_{1 \le i \le n} A_{ji} x_i \ge b_j
namespace simplex {
// maximize c^Tx under Ax <= B
                                                                    4. If x_i has no lower bound, replace x_i with x_i - x_i'
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
                                                                       Geometry
using VD = vector<double>;
using VVD = vector<vector<double>>;
                                                                  6.1 Basic Geometry
const double eps = 1e-9;
const double inf = 1e+9;
                                                                  #define IM imag
                                                                  #define RE real
int n, m;
VVD d;
                                                                  using lld = int64_t;
vector<int> p, q;
                                                                  using llf = long double;
                                                                  using PT = std::complex<lld>;
void pivot(int r, int s) {
 double inv = 1.0 / d[r][s];
                                                                  using PTF = std::complex<llf>;
                                                                  auto toPTF(PT p) { return PTF{RE(p), IM(p)}; } int sgn(lld x) { return (x > 0) - (x < 0); }
 for (int i = 0; i < m + 2; ++i)
for (int j = 0; j < n + 2; ++j)
   if (i != r && j != s)
                                                                  11d dot(PT a, PT b) { return RE(conj(a) * b); }
    d[i][j] = d[r][j] * d[i][s] * inv;
                                                                  11d cross(PT a, PT b) { return \ IM(conj(a) * b);  }
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;</pre>
                                                                  int ori(PT a, PT b, PT c) {
 for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
                                                                  return sgn(cross(b - a, c - a));
 d[r][s] = inv; swap(p[r], q[s]);
                                                                  bool operator<(const PT &a, const PT &b) {</pre>
                                                                  return RE(a) != RE(b) ? RE(a) < RE(b) : IM(a) < IM(b);
bool phase(int z) {
 int x = m + z;
                                                                  int quad(PT p) {
  return (IM(p) == 0) // use sgn for PTF
 while (true) {
  int s = -1
  for (int i = 0; i <= n; ++i) {
                                                                    ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ? 0 : 2);
   if (!z && q[i] == -1) continue;
   if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
                                                                  int argCmp(PT a, PT b) {
                                                                   // -1 / 0 / 1 <-> < / == / > (atan2)
  if (d[x][s] > -eps) return true;
                                                                   int qa = quad(a), qb = quad(b);
                                                                   if (qa != qb) return sgn(qa - qb);
  for (int i = 0; i < m; ++i) {
                                                                   return sgn(cross(b, a));
```

template <typename V> llf area(const V & pt) {

11d ret = 0;

REP(i,n) REP(j,n) flag[i][j] = 0;

```
for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
                                                               vector<Face> now;
  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
                                                               now.emplace_back(0,1,2);
 return ret / 2.0;
                                                               now.emplace_back(2,1,0);
                                                               for (int i=3; i<n; i++){
PT rot90(PT p) { return PT{-IM(p), RE(p)}; }
                                                                ftop++; vector<Face> next;
PTF project(PTF p, PTF q) { // p onto q
return dot(p, q) * q / dot(q, q);
                                                                REP(j, SZ(now)) {
                                                                 Face& f=now[j]; int ff = 0;
                                                                 ld d=(pt[i]-pt[f.a]).dot(
11f FMOD(11f x) {
                                                                   ver(pt[f.a], pt[f.b], pt[f.c]));
 if (x < -PI) x += PI * 2;
                                                                 if (d <= 0) next.push_back(f);</pre>
 if (x > PI) x -= PI * 2;
                                                                 if (d > 0) ff=ftop;
                                                                 else if (d < 0) ff=-ftop;</pre>
 return x;
                                                                 flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
                                                                REP(j, SZ(now)) {
      Segment & Line Intersection
                                                                 Face& f=now[j]
struct Segment {
                                                                 if (flag[f.a][f.b] > 0 &&
 PT st, dir; // represent st + t*dir for 0<=t<=1
                                                                   flag[f.a][f.b] != flag[f.b][f.a])
 Segment(PT s, PT e) : st(s), dir(e - s) \{\}
                                                                  next.emplace_back(f.a,f.b,i);
 static bool valid(lld p, lld q) {
                                                                 if (flag[f.b][f.c] > 0 &&
  // is there t s.t. 0 <= t <= 1 && qt == p ?
                                                                   flag[f.b][f.c] != flag[f.c][f.b])
  if (q < 0) q = -q, p = -p;
                                                                  next.emplace_back(f.b,f.c,i);
  return 0 <= p && p <= q;
                                                                 if (flag[f.c][f.a] > 0 &&
  flag[f.c][f.a] != flag[f.a][f.c])
                                                                  next.emplace_back(f.c,f.a,i);
bool isInter(Segment A, PT P) {
if (A.dir == PT(0)) return P == A.st;
return cross(P - A.st, A.dir) == 0 &&
                                                                now=next;
  Segment::valid(dot(P - A.st, A.dir), norm(A.dir));
                                                               return now;
template <typename U, typename V>
bool isInter(U A, V B) {
  if (cross(A.dir, B.dir) == 0)
                                                              6.5 2D Farthest Pair
                                                              // stk is from convex hull
  return // handle parallel yourself
                                                              n = (int)(stk.size());
   isInter(A, B.st) || isInter(A, B.st+B.dir) ||
isInter(B, A.st) || isInter(B, A.st+A.dir);
                                                              int pos = 1, ans = 0; stk.push_back(stk[0]);
                                                              for(int i=0;i<n;i++) {</pre>
 PT D = B.st - A.st;
                                                               while(abs(cross(stk[i+1]-stk[i],
 11d C = cross(A.dir, B.dir);
                                                                 stk[(pos+1)%n]-stk[i])) >
 return U::valid(cross(D, A.dir), C) &&
                                                                 abs(cross(stk[i+1]-stk[i],
   V::valid(cross(D, B.dir), C);
                                                                 stk[pos]-stk[i]))) pos = (pos+1)%n;
                                                               ans = max({ans, dis(stk[i], stk[pos]),
struct Line
                                                                dis(stk[i+1], stk[pos])});
PT st, ed, dir;
Line (PT s, PT e)
  : st(s), ed(e), dir(e - s) {}
                                                              6.6 kD Closest Pair (3D ver.)
                                                              11f solve(vector<P> v) {
PTF intersect(const Line &A, const Line &B) {
                                                               shuffle(v.begin(), v.end(), mt19937());
 11f t = cross(B.st - A.st, B.dir) /
 llf(cross(A.dir, B.dir))
                                                               unordered_map<11d, unordered_map<11d,</pre>
                                                                unordered_map<lld, int>>> m;
 return toPTF(A.st) + PTF(t) * toPTF(A.dir);
                                                               llf d = dis(v[0], v[1]);
                                                               auto Idx = [&d] (11f x) -> 11d {
                                                                return round(x * 2 / d) + 0.1;
6.3 2D Convex Hull
                                                               auto rebuild_m = [&m, &v, &Idx](int k) {
void make_hull(vector<pll> &dots) { // n=1 => ans = {}
                                                                m.clear();
 sort(dots.begin(), dots.end());
                                                                for (int i = 0; i < k; ++i)
 vector<pll> ans(1, dots[0]);
                                                                 m[Idx(v[i].x)][Idx(v[i].y)]
 for (int ct = 0; ct < 2; ++ct, reverse(ALL(dots)))</pre>
                                                                  [Idx(v[i].z)] = i;
  for (int i = 1, t = SZ(ans); i < SZ(dots); i++) {</pre>
                                                               }; rebuild_m(2);
   while (SZ(ans) > t && ori(
                                                               for (size_t i = 2; i < v.size(); ++i) {</pre>
     ans[SZ(ans) - 2], ans.back(), dots[i]) <= 0
                                                                const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
    ans.pop_back();
                                                                   kz = Idx(v[i].z); bool found = false;
   ans.pb(dots[i]);
                                                                for (int dx = -2; dx <= 2; ++dx) {
                                                                 const 11d nx = dx + kx;
 ans.pop_back(), ans.swap(dots);
                                                                 if (m.find(nx) == m.end()) continue;
                                                                 auto\& mm = m[nx];
                                                                 for (int dy = -2; dy <= 2; ++dy) {
      3D Convex Hull
                                                                  const 11d ny = dy + ky;
// return the faces with pt indexes
                                                                  if (mm.find(ny) == mm.end()) continue;
                                                                  auto& mmm = mm[ny];
int flag[MXN][MXN];
struct Point{
                                                                  for (int dz = -2; dz <= 2; ++dz) {
 ld x,y,z;
                                                                   const 11d nz = dz + kz;
 Point operator * (const ld &b) const {
                                                                   if (mmm.find(nz) == mmm.end()) continue;
  return (Point){x*b,y*b,z*b};}
                                                                   const int p = mmm[nz];
 Point operator * (const Point &b) const {
                                                                   if (dis(v[p], v[i]) < d) {</pre>
  return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
                                                                    d = dis(v[p], v[i]);
                                                                    found = true;
                                                                   }
Point ver(Point a, Point b, Point c) {
return (b - a) * (c - a);}
vector<Face> convex_hull_3D(const vector<Point> pt) {
 int n = SZ(pt), ftop = 0;
                                                                if (found) rebuild_m(i + 1);
```

else m[kx][ky][kz] = i;

```
PTF dir = B.o - A.o; llf d2 = norm(dir);
                                                               if (norm(A.r - B.r) >= d2) // norm(x) := |x|^2
 return d;
                                                               if (A.r < B.r) return {-PI, PI}; // A in B</pre>
                                                                else return {}; // B in A
6.7
      Simulated Annealing
                                                               if (norm(A.r + B.r) <= d2) return {};</pre>
                                                               11f dis = abs(dir), theta = arg(dir);
11f anneal() {
                                                               11f phi = acos((A.r * A.r + d2 - B.r * B.r) /
 mt19937 rnd_engine( seed );
                                                                 (2 * A.r * dis));
 uniform_real_distribution< llf > rnd( 0, 1 );
                                                               11f L = FMOD(theta - phi), R = FMOD(theta + phi);
 const 11f dT = 0.001;
                                                               return { L, R };
 // Argument p
 11f S_cur = calc( p ), S_best = S_cur;
for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
                                                             vector<PTF> intersectPoint(Circle a, Circle b) {
  // Modify p to p_prime
                                                              llf d = abs(a.o - b.o);
  const llf S_prime = calc( p_prime );
                                                               if (d >= b.r+a.r || d <= abs(b.r-a.r)) return {};</pre>
  const llf delta_c = S_prime - S_cur;
  llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
                                                               11f dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
                                                               PTF dir = (a.o - b.o) / d;
  if ( rnd( rnd_engine ) <= prob )</pre>
                                                               PTF u = dir*d1 + b.o;
 S_cur = S_prime, p = p_prime;
if ( S_prime < S_best ) // find min</pre>
                                                              PTF v = rot90(dir) * sqrt(max<llf>(0, b.r*b.r-d1*d1));
                                                               return \{u + v, u - v\};
   S_best = S_prime, p_best = p_prime;
 return S_best;
                                                             6.11 Intersection of line and Circle
                                                             vector<PTF> line_interCircle(const PTF &p1,
6.8 Half Plane Intersection
                                                                const PTF &p2, const PTF &c, const double r)
                                                               PTF ft = p1 + project(c-p1, p2-p1), vec = p2-p1;
// cross(pt-line.st, line.dir)<=0 <-> pt in half plane
bool operator<(const Line &lhs, const Line &rhs) {
                                                               llf dis = abs(c - ft);
                                                               if (abs(dis - r) < eps) return {ft};</pre>
  if (int cmp = argCmp(lhs.dir, rhs.dir))
                                                              if (dis > r) return {};
vec = vec * sqrt(r * r - dis * dis) / abs(vec);
    return cmp == -1;
  return ori(lhs.st, lhs.ed, rhs.st) < 0;
                                                               return {ft + vec, ft - vec};
// intersect function is in "Segment Intersect"
                                                             6.12 Intersection of Polygon and Circle
11f HPI(vector<Line> &lines) {
  sort(lines.begin(), lines.end());
                                                             // Divides into multiple triangle, and sum up
  deque<Line> que;
                                                             // test by HDU2892
  deque<PTF> pt;
                                                             11f _area(PTF pa, PTF pb, llf r)
  que.push_back(lines[0]);
                                                              if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
  for (int i = 1; i < (int)lines.size(); i++) {</pre>
                                                               if (abs(pb) < eps) return 0;</pre>
    if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
                                                               llf S, h, theta;
     continue;
                                                               11f a = abs(pb), b = abs(pa), c = abs(pb - pa);
#define POP(L, R) \
                                                               11f cosB = dot(pb, pb - pa) / a / c, B = acos(cosB);
    while (pt.size() > 0 \
                                                               11f cosC = dot(pa, pb) / a / b, C = acos(cosC);
      && ori(L.st, L.ed, pt.back()) < 0) \
                                                               if (a > r) {
    pt.pop_back(), que.pop_back(); \
while (pt.size() > 0 \
                                                               S = (C / 2) * r * r;
                                                                h = a * b * sin(C) / c;
      && ori(R.st, R.ed, pt.front()) < 0) \
                                                               if (h < r && B < PI / 2)
      pt.pop_front(), que.pop_front();
                                                                 S = (acos(h / r) * r * r - h * sqrt(r*r - h*h));
    POP(lines[i], lines[i]);
                                                               } else if (b > r) {
  theta = PI - B - asin(sin(B) / r * a);
    pt.push_back(intersect(que.back(), lines[i]));
    que.push_back(lines[i]);
                                                                S = 0.5 * a*r * sin(theta) + (C - theta) / 2 * r*r;
                                                               } else
  POP(que.front(), que.back())
                                                               S = 0.5 * sin(C) * a * b;
  if (que.size() <= 1 ||</pre>
                                                               return S;
    argCmp(que.front().dir, que.back().dir) == 0)
                                                             11f area_poly_circle(const vector<PTF> &poly,
  pt.push_back(intersect(que.front(), que.back()));
                                                               const PTF &0, const llf r) {
  return area(pt);
                                                               11f S = 0:
                                                               for (int i = 0, N = poly.size(); i < N; ++i)</pre>
                                                                S += _area(poly[i] - 0, poly[(i + 1) % N] - 0, r) *
6.9 Minkowski Sum
                                                                   ori(0, poly[i], poly[(i + 1) % N]);
vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
                                                               return fabs(S);
hull(A), hull(B);
 vector<pll> C(1, A[0] + B[0]), s1, s2;
 for(int i = 0; i < SZ(A); ++i)
s1.pb(A[(i + 1) % SZ(A)] - A[i]);
                                                             6.13 Point & Hulls Tangent
                                                             #define above(P, Vi, Vj) (ori(P, Vi, Vj) > 0) // true
 for(int i = 0; i < SZ(B); i++)
                                                                  if Vi is above Vj
  s2.pb(B[(i + 1) % SZ(B)] - B[i]);
                                                             #define below(P, Vi, Vj) (ori(P, Vi, Vj) < 0) // true</pre>
 for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
                                                                  if Vi is below Vj
  if (p2 >= SZ(B)
                                                             // Rtangent_PointPolyC(): binary search for convex
    || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
                                                                  polygon right tangent
   C.pb(C.back() + s1[p1++]);
                                                                  Input: P = a 2D point (exterior to the polygon)
  else
                                                             //
                                                                      n = number of polygon vertices
   C.pb(C.back() + s2[p2++]);
                                                             //
                                                                      V = array of vertices for a 2D convex polygon
 return hull(C), C;
                                                                  with V[n] = V[0]
                                                                  Return: index "i" of rightmost tangent point V[i]
                                                             int Rtangent_PointPolyC(PT P, int n, PT *V) {
6.10 Circle Class
                                                              int a, b, c;
struct Circle { PTF o; llf r; };
                                                              int upA, dnC;
vector<llf> intersectAngle(Circle A, Circle B) {
                                                              if (below(P, V[1], V[0]) && !above(P, V[n - 1], V[0]))
```

```
return 0;
for (a = 0, b = n;;) {
 c = (a + b) / 2
 dnC = below(P, V[c + 1], V[c]);
  if (dnC && !above(P, V[c - 1], V[c]))
  upA = above(P, V[a + 1], V[a]);
  if (upA) {
  if (dnC) {
   b = c;
   } else {
    if (above(P, V[a], V[c]))
     b = c;
    else
     a = c:
  } else {
   if (!dnC) {
   a = c;
   } else {
   if (below(P, V[a], V[c]))
    b = c;
    else
    a = c;
   }
  }
// Ltangent_PointPolyC(): binary search for convex
    polygon left tangent
    Input: P = a 2D point (exterior to the polygon)
        n = number of polygon vertices
//
        V = array of vertices for a 2D convex polygon
11
    with V[n]=V[0]
   Return: index "i" of leftmost tangent point V[i]
int Ltangent_PointPolyC(PT P, int n, PT *V) {
int a, b, c;
int dnA, dnC;
if (above(P, V[n - 1], V[0]) && !below(P, V[1], V[0]))
for (a = 0, b = n;;) {
 c = (a + b) / 2;
 dnC = below(P, V[c + 1], V[c]);
  if (above(P, V[c - 1], V[c]) && !dnC)
   return c:
  dnA = below(P, V[a + 1], V[a]);
  if (dnA) {
  if (!dnC) {
   b = c;
   } else {
    if (below(P, V[a], V[c]))
     b = c;
    else
     a = c:
  } else {
   if (dnC) {
   a = c;
   } else {
    if (above(P, V[a], V[c]))
     b = c;
   else
     a = c:
  }
}
}
      Convex Hulls Tangent
```

```
// RLtangent_PolyPolyC(): get the RL tangent between
    two convex polygons
// Input: m = number of vertices in polygon 1
// V = array of vertices for convex polygon 1 with
    V[m]=V[0]
// n = number of vertices in polygon 2
// W = array of vertices for convex polygon 2 with
    W[n]=W[0]
```

```
Output: *t1 = index of tangent point V[t1] for
    polygon 1
        *t2 = index of tangent point W[t2] for polygon
11
void RLtangent_PolyPolyC(int m, PT *V, int n, PT *W,
    int *t1, int *t2) {
 int ix1, ix2; // search indices for polygons 1 and 2
 // first get the initial vertex on each polygon
 ix1 = Rtangent_PointPolyC(W[0], m, V); // right
    tangent from W[\theta] to V
 ix2 = Ltangent_PointPolyC(V[ix1], n, W); // left
    tangent from V[ix1] to W
 // ping-pong linear search until it stabilizes
 int done = false; // flag when done
 while (done == false) {
  done = true; // assume done until..
  while (ori(W[ix2], V[ix1], V[ix1 + 1]) <= 0) {</pre>
   ++ix1; // get Rtangent from W[ix2] to V
  while (ori(V[ix1], W[ix2], W[ix2 - 1]) >= 0) {
            // get Ltangent from V[ix1] to W
   done = false; // not done if had to adjust this
  }
 *t1 = ix1;
 *t2 = ix2;
 return;
```

6.15 Tangent line of Two Circle

```
vector<Line>
tanline(const Circle &c1, const Circle &c2, int sign1){
 // sign1 = 1 for outer tang, -1 for inter tang
 vector<Line> ret;
 if (norm(c1.o - c2.o) < eps) return ret;</pre>
 11f d = abs(c1.o - c2.o);
 PTF v = (c2.o - c1.o) / d;
11f c = (c1.r - sign1 * c2.r) / d;
 if (c * c > 1) return ret;
 llf h = sqrt(max<llf>(0, 1 - c * c));
 for (int sign2 : {1, -1}) {
  PTF n = c * v + sign2 * h * rot90(v);
  PTF p1 = c1.o + n \star c1.r;
PTF p2 = c2.o + n \star (c2.r \star sign1);
  if (norm(p2 - p1) < eps)
   p2 = p1 + rot90(c2.o - c1.o);
  ret.push_back({p1, p2});
 return ret;
}
```

6.16 Minimum Covering Circle

template<typename P>

```
Circle getCircum(const P &a, const P &b, const P &c){
 Real a1 = a.x-b.x, b1 = a.y-b.y;
 Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
 Real a2 = a.x-c.x, b2 = a.y-c.y;
 Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
 Circle cc;
 cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
 cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
 cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
 return cc;
}
template<typename P>
Circle MinCircleCover(const vector<P>& pts){
 random_shuffle(pts.begin(), pts.end());
 Circle c = { pts[0], 0 };
 for(int i=0;i<(int)pts.size();i++){</pre>
  if (dist(pts[i], c.o) <= c.r) continue;</pre>
  c = { pts[i], 0 };
for (int j = 0; j < i; j++) {
  if(dist(pts[j], c.o) <= c.r) continue;</pre>
   c.o = (pts[i] + pts[j]) / 2;
   c.r = dist(pts[i], c.o)
   for (int k = 0; k < j; k++) {
```

if (dist(pts[k], c.o) <= c.r) continue;</pre>

c = getCircum(pts[i], pts[j], pts[k]);

```
6.18
                                                                        Rotating Sweep Line
  }
                                                                 void rotatingSweepLine(pair<int, int> a[], int n) {
 return c;
                                                                   vector<pair<int, int>> 1;
                                                                   1.reserve(n * (n - 1) / 2);
                                                                   for (int i = 0; i < n; ++i)
for (int j = i + 1; j < n; ++j)
       KDTree (Nearest Point)
const int MXN = 100005;
                                                                     l.emplace_back(i, j)
struct KDTree {
                                                                   sort(1.begin(), 1.end(), [&a](auto &u, auto &v){
struct Node {
                                                                    11d udx = a[u.first].first - a[u.second].first;
  int x,y,x1,y1,x2,y2;
                                                                    11d udy = a[u.first].second - a[u.second].second;
 int id,f;
Node *L, *R;
                                                                    1ld vdx = a[v.first].first - a[v.second].first;
                                                                    11d vdy = a[v.first].second - a[v.second].second;
 } tree[MXN], *root;
                                                                    if (udx == 0 or vdx == 0) return not udx == 0;
                                                                    int s = sgn(udx * vdx);
 LL dis2(int x1, int y1, int x2, int y2) {
                                                                    return udy * vdx * s < vdy * udx * s;
 LL dx = x1-x2, dy = y1-y2;
  return dx*dx+dy*dy;
                                                                   vector<int> idx(n), p(n);
                                                                   iota(idx.begin(), idx.end(), 0);
sort(idx.begin(), idx.end(), [&a](int i, int j){
 static bool cmpx(Node& a, Node& b){return a.x<b.x;}
static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
                                                                   return a[i] < a[j]; });
 void init(vector<pair<int,int>> ip) {
                                                                   for (int i = 0; i < n; ++i) p[idx[i]] = i;</pre>
  n = ip.size();
                                                                   for (auto [i, j]: 1) {
  for (int i=0; i<n; i++) {</pre>
                                                                   // do here
   tree[i].id = i;
                                                                    swap(p[i], p[j]);
   tree[i].x = ip[i].first;
                                                                   idx[p[i]] = i, idx[p[j]] = j;
   tree[i].y = ip[i].second;
                                                                 }
  root = build_tree(0, n-1, 0);
                                                                 6.19
                                                                        Circle Cover
 Node* build_tree(int L, int R, int d) {
  if (L>R) return nullptr
                                                                 const int N = 1021;
  int M = (L+R)/2; tree[M].f = d%2;
                                                                  struct CircleCover {
  nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
                                                                   int C;
  tree[M].x1 = tree[M].x2 = tree[M].x;
                                                                   Cir c[N]
  tree[M].y1 = tree[M].y2 = tree[M].y
                                                                   bool g[N][N], overlap[N][N];
                                                                   // Area[i] : area covered by at least i circles
double Area[ N ];
  tree[M].L = build_tree(L, M-1, d+1);
  if (tree[M].L) {
   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
                                                                   void init(int _C){ C = _C;}
   tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
                                                                   struct Teve {
                                                                   PTF p; double ang; int add;
   tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
                                                                    Teve() {}
                                                                    Teve(PTF _a, double _b, int _c):p(_a), ang(_b), add(
  tree[M].R = build_tree(M+1, R, d+1);
  if (tree[M].R) {
                                                                    bool operator<(const Teve &a)const
   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
                                                                    {return ang < a.ang;}</pre>
                                                                   }eve[N * 2];
   tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
                                                                   // strict: x = 0, otherwise x = -1
bool disjuct(Cir &a, Cir &b, int x)
   tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
                                                                   {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
                                                                   bool contain(Cir &a, Cir &b, int x)
  return tree+M;
                                                                   \{return sign(a.R - b.R - abs(a.0 - b.0)) > x;\}
 int touch(Node* r, int x, int y, LL d2){
                                                                   bool contain(int i, int j) {
                                                                   /* c[j] is non-strictly in c[i]. */
 LL dis = sqrt(d2)+1;
  if (x<r->x1-dis || x>r->x2+dis ||
                                                                    return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c[i].R) \mid c[i].R - c[i].R)
    y<r->y1-dis || y>r->y2+dis)
                                                                      [j].R) == 0 \&\& i < j)) \&\& contain(c[i], c[j], -1);
   return 0;
                                                                   void solve(){
  return 1;
                                                                    fill_n(Area, C + 2, 0);
                                                                    for(int i = 0; i < C; ++i)
for(int j = 0; j < C; ++j)</pre>
 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
  if (!r || !touch(r, x, y, md2)) return;
  LL d2 = dis2(r\rightarrow x, r\rightarrow y, x, y)
                                                                      overlap[i][j] = contain(i, j);
  if (d2 < md2 \mid \mid (d2 == md2 \&\& mID < r->id)) {
                                                                    for(int i = 0; i < C; ++i)
  mID = r -> id;
                                                                     for(int j = 0; j < C; ++j)
                                                                      g[i][j] = !(overlap[i][j] || overlap[j][i] ||
   md2 = d2;
                                                                        disjuct(c[i], c[j], -1));
  // search order depends on split dim
                                                                    for(int i = 0; i < C; ++i){</pre>
  if ((r->f == 0 && x < r->x) ||
                                                                     int E = 0, cnt = 1;
    (r->f == 1 \&\& y < r->y)) {
                                                                     for(int j = 0; j < C; ++j)
                                                                      if(j != i && overlap[j][i])
   nearest(r\rightarrow L, x, y, mID, md2);
   nearest(r->R, x, y, mID, md2);
                                                                       ++cnt:
                                                                     for(int j = 0; j < C; ++j)
if(i != j && g[i][j]) {</pre>
  } else {
   nearest(r->R, x, y, mID, md2);
   nearest(r->L, x, y, mID, md2);
                                                                       auto IP = intersectPoint(c[i], c[j]);
                                                                       PTF aa = IP[0], bb = IP[1];
                                                                       llf A = arg(aa-c[i].0), B = arg(bb-c[i].0);
 int query(int x, int y) {
                                                                       eve[E++] = Teve(bb,B,1), eve[E++]=Teve(aa,A,-1);
                                                                       if(B > A) ++cnt;
  int id = 1029384756;
  LL d2 = 102938475612345678LL;
  nearest(root, x, y, id, d2);
                                                                     if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
  return id;
                                                                     else{
                                                                      sort(eve, eve + E);
} tree;
                                                                      eve[E] = eve[0];
```

```
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    for(int j = 0; j < E; ++j){}
     cnt += eve[j].add;
     Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
     double theta = eve[j + 1].ang - eve[j].ang;
     if (theta < 0) theta += 2. * pi;</pre>
     Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
7
    Stringology
     Suffix Array
namespace sfx {
bool _t[maxn * 2];
```

```
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], _c[maxn * 2];
int x[maxn], _p[maxn], _q[maxn * 2];
// sa[i]: sa[i]-th suffix is the
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *a, int *c, int n, int z) {
 memset(a, 0, sizeof(int) * n);
 memcpy(x, c, sizeof(int) * z);
void induce(int *a,int *c,int *s,bool *t,int n,int z){
 memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
  if (a[i] && !t[a[i] - 1])
   a[x[s[a[i] - 1]]++] = a[i] - 1;
 memcpy(x, c, sizeof(int) * z);
 for (int i = n - 1; i >= 0; --i)
if (a[i] && t[a[i] - 1])
   a[--x[s[a[i] - 1]]] = a[i] - 1;
void sais(int *s, int *a, int *p, int *q,
bool *t, int *c, int n, int z) {
 bool uniq = t[n - 1] = true;
 int nn=0, nmxz=-1, *nsa = a+n, *ns=s+n, last=-1;
 memset(c, 0, sizeof(int) * z);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
 for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
 if (uniq) {
  for (int i = 0; i < n; ++i) a[--c[s[i]]] = i;
  return;
 for (int i = n - 2; i \ge 0; --i)
  t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 pre(a, c, n, z);
for (int i = 1; i <= n - 1; ++i)</pre>
  if (t[i] && !t[i - 1])
   a[--x[s[i]]] = p[q[i] = nn++] = i;
 induce(a, c, s, t, n, z);
for (int i = 0; i < n; ++i)
  if (a[i] && t[a[i]] && !t[a[i] - 1]) {
  bool neq = last < 0 || \</pre>
   memcmp(s + a[i], s + last,
(p[q[a[i]] + 1] - a[i]) * sizeof(int));
  ns[q[last = a[i]]] = nmxz += neq;
 }}
 sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
 pre(a, c, n, z);
for (int i = nn - 1; i >= 0; --i)
  a[--x[s[p[nsa[i]]]]] = p[nsa[i]];
 induce(a, c, s, t, n, z);
void build(const string &s) {
 const int n = int(s.size());
 for (int i = 0; i < n; ++i) _s[i] = s[i];
 _s[n] = 0; // s shouldn't contain 0
 sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
for(int i = 0; i < n; ++i) rev[sa[i] = sa[i+1]] = i;</pre>
 int ind = hi[0] = 0;
 for (int i = 0; i < n; ++i) {
  if (!rev[i]) {</pre>
   ind = 0;
   continue;
  while (i + ind < n && \</pre>
```

```
s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  hi[rev[i]] = ind ? ind-- : 0;
}}
7.2 Suffix Automaton
 struct node {
  int ch[K], len, fail, cnt, indeg;
```

```
struct SuffixAutomaton {
  node(int L = 0) : ch{}, len(L), fail(0), cnt(0),
    indea(0) {}
 } st[N];
 int root, last, tot;
 void extend(int c) {
  int cur = ++tot;
  st[cur] = node(st[last].len + 1);
  while (last && !st[last].ch[c]) {
    st[last].ch[c] = cur;
    last = st[last].fail;
  if (!last) {
    st[cur].fail = root;
  } else {
    int q = st[last].ch[c];
    if (st[q].len == st[last].len + 1) {
      st[cur].fail = q;
    } else {
      int clone = ++tot;
      st[clone] = st[q];
      st[clone].len = st[last].len + 1;
      st[st[cur].fail = st[q].fail = clone].cnt = 0;
      while (last && st[last].ch[c] == q) {
        st[last].ch[c] = clone;
        last = st[last].fail;
    }
  st[last = cur].cnt += 1;
 void init(const char* s) {
  root = last = tot = 1;
  st[root] = node(0);
  for (char c; c = *s; ++s) extend(c - 'a');
 int q[N];
 void dp() {
  for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
  int head = 0, tail = 0;
  for (int i = 1; i <= tot; i++)</pre>
    if (st[i].indeg == 0) q[tail++] = i;
  while (head != tail) {
    int now = q[head++];
    if (int f = st[now].fail) {
      st[f].cnt += st[now].cnt;
      if (--st[f].indeg == 0) q[tail++] = f;
  }
 int run(const char* s) {
  int now = root;
  for (char_c; c = *s; ++s) {
    if (!st[now].ch[c -= 'a']) return 0;
    now = st[now].ch[c];
  return st[now].cnt;
 }
} SAM;
```

7.3 Z value

```
vector<int> Zalgo(const string &s) {
 vector<int> z(s.size(), s.size());
 for (int i = 1, 1 = 0, r = 0; i < z[0]; ++i) {
  int j = clamp(r - i, 0, z[i - 1]);</pre>
  for (; i + j < z[0] \text{ and } s[i + j] == s[j]; ++j);
  if (i + (z[i] = j) > r) r = i + z[1 = i];
 return z;
```

7.4 Manacher

```
int z[maxn];
int manacher(const string& s) {
    string t = ".";
    for(char c: s) t += c, t += '.';
    int l = 0, r = 0, ans = 0;
    for (int i = 1; i < t.length(); ++i) {
        z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
        while (i - z[i] >= 0 && i + z[i] < t.length()) {
        if(t[i - z[i]] == t[i + z[i]]) ++z[i];
        else break;
    }
    if (i + z[i] > r) r = i + z[i], l = i;
}
for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
    return ans;
}</pre>
```

7.5 Lexico Smallest Rotation

```
string mcp(string s) {
  int n = s.length();
  s += s; int i = 0, j = 1;
  while (i < n && j < n) {
    int k = 0;
    while (k < n && s[i + k] == s[j + k]) k++;
    ((s[i + k] <= s[j + k]) ? j : i) += k + 1;
    j += (i == j);
  }
  return s.substr(i < n ? i : j, n);
}</pre>
```

7.6 Main Lorentz

```
vector<tuple<tuple<size_t, size_t, int, int>>> reps;
void find_repetitions(const string &s, int shift = 0) {
if (s.size() <= 1)
  return
 const size_t nu = s.size() / 2, nv = s.size() - nu;
string u = s.substr(0, nu), v = s.substr(nu);
string ru(u.rbegin(), u.rend());
string rv(v.rbegin(), v.rend());
find_repetitions(u, shift);
find_repetitions(v, shift + nu);
auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u),
    z3 = Zalgo(ru + '#' + rv), z4 = Zalgo(v);
 for (size_t cntr = 0; cntr < s.size(); cntr++) {</pre>
 size_t 1; int k1, k2;
if (cntr < nu) {</pre>
   1 = nu - cntr;
   k1 = 1 < z1.size() ? z1[1] : 0;
   k2 = n + 1 - 1 < z2.size() ? z2[n + 1 - 1] : 0;
  } else {
   1 = cntr - nu + 1;
   k1 = n + 1 - 1 < z3.size() ? z3[n + 1 - 1] : 0;
   k2 = 1 < z4.size() ? z4[1] : 0;
  if (k1 + k2 >= 1)
   reps.emplace_back(cntr, 1, k1, k2);
```

7.7 BWT

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a'
vector<int> v[ SIGMA ];
void BWT(char* ori, char* res){
 // make ori -> ori + ori
 // then build suffix array
void iBWT(char* ori, char* res){
 for( int i = 0 ; i < SIGMA ; i ++ )
  v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )</pre>
  v[ ori[i] - BASE ].push_back( i );
  vector<int> a:
 for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )</pre>
  for( auto j : v[ i ] ){
   a.push_back( j );
    ori[ ptr ++ ] = BASE + i;
```

```
for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
   ptr = a[ ptr ];
  res[ len ] = 0;
 }
} bwt;
7.8 Palindromic Tree
struct palindromic_tree{
 struct node{
  int next[26],f,len;
  int cnt, num, st, ed;
  node(int 1=0):f(0),len(1),cnt(0),num(0) {
   memset(next, 0, sizeof(next)); }
 vector<node> st;
 vector<char> s;
 int last,n;
 void init(){
  st.clear();s.clear();last=1; n=0;
  st.push_back(0);st.push_back(-1);
  st[0].f=1;s.push_back(-1); }
 int getFail(int x){
  while(s[n-st[x].len-1]!=s[n])x=st[x].f;
  return x;}
 void add(int c){
  s.push_back(c-='a'); ++n;
  int cur=getFail(last);
  if(!st[cur].next[c]){
   int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
   st[now].num=st[st[now].f].num+1;
  last=st[cur].next[c];
  ++st[last].cnt;}
 void dpcnt() {
  for (int i=st.size()-1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
 int size(){ return st.size()-2;}
} pt;
int main() {
 string s; cin >> s; pt.init();
 for (int i=0; i<SZ(s); i++) {
  int prvsz = pt.size(); pt.add(s[i]);</pre>
  if (prvsz != pt.size()) {
   int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [l,r]: s.substr(l, r-l+1)
 return 0;
8
     Misc
```

8.1 Theorems

8.1.1 Sherman-Morrison formula

$$(A + uv^{\mathsf{T}})^{-1} = A^{-1} - \frac{A^{-1}uv^{\mathsf{T}}A^{-1}}{1+v^{\mathsf{T}}A^{-1}u}$$

8.1.2 Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

8.1.3 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.4 Cayley's Formula

- Given a degree sequence d_1,d_2,\dots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.5 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \le k \le n$.

8.1.6 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.7 Euler's planar graph formula

V - E + F = C + 1, $E \le 3V - 6$ (?)

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have A #{lattice points in the interior} + $\frac{\#\{\text{lattice points on the boundary}\}}{2} - 1$

8.1.9 Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2),$ find maximum $S\in I_1\cap I_2.$ For each iteration, build the directed graph and find a shortest path from s to t.

- $s \rightarrow x : S \sqcup \{x\} \in I_1$
- $x \rightarrow t : S \sqcup \{x\} \in I_2$
- $y \to x: S \setminus \{y\} \sqcup \{x\} \in I_1$ (y is in the unique circuit of $S \sqcup \{x\}$)
- $x \to y : S \setminus \{y\} \sqcup \{x\} \in I_2$ (y is in the unique circuit of $S \sqcup \{x\}$)

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|$. In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

8.2 Bitset LCS

```
scanf("%d%d", &n, &m), u = n / 64 + 1;
for (int i = 1, c; i <= n; i++)
    scanf("%d", &c), p[c].set(i);
for (int i = 1, c; i <= m; i++) {
    scanf("%d", &c), (g = f) |= p[c];
    f.shiftLeftByOne(), f.set(0);
    ((f = g - f) ^= g) &= g;
}
printf("%d\n", f.count());</pre>
```

8.3 Convex 1D/1D DP

```
struct segment {
int i, 1, r
segment() {}
segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
dp[0] = 0;
deque<segment> dq; dq.push_back(segment(0, 1, n));
for (int i = 1; i <= n; ++i) {
 dp[i] = f(dq.front().i, i);
 while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
 dq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
 while (dq.size() &&
   f(i, dq.back().1) < f(dq.back().i, dq.back().1))
    dq.pop_back();
  if (dq.size())
  int d = 1 << 20, c = dq.back().1;</pre>
  while (d \gg 1) if (c + d \ll dq.back().r)
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
 if (seg.1 <= n) dq.push_back(seg);</pre>
```

8.4 ConvexHull Optimization

```
struct L
 mutable int64_t a, b, p;
 bool operator<(const L &r) const { return a < r.a; }</pre>
 bool operator<(int64_t x) const { return p < x; }</pre>
struct DynamicHull : multiset<L, less<>> {
 static const int64_t kInf = 1e18;
 bool Isect(iterator x, iterator y)
  auto Div = [](int64_t a, int64_t b) {
    return a / b - ((a ^ b) < 0 && a % b); }
  if (y == end()) { x->p = kInf; return false; }
if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(int64_t a, int64_t b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
while (Isect(y, z)) z = erase(z);
  if (x!=begin()&&Isect(--x,y)) Isect(x, y=erase(y));
  while ((y = x) != begin() \&\& (--x)->p >= y->p)
   Isect(x, erase(y));
 int64_t Query(int64_t x) {
  auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
};
```

8.5 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

8.6 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
 dfn[u]=low[u]=++dfs_idx;
 for(int i=0;i<(int)init_g[u].size();i++){</pre>
  int v=init_g[u][i];
  if(v==par[u]) continue;
  if(!dfn[v]){
   par[v]=u;
   tarjan(v)
   low[u]=min(low[u],low[v]);
   if(dfn[u]<low[v]){</pre>
    g[u].push_back(v)
    g[v].push_back(u);
  }else{
   low[u]=min(low[u],dfn[v]);
   if(dfn[v]<dfn[u]){</pre>
    int temp_v=u;
    bcc_id++;
    while(temp_v!=v){
     g[bcc_id+n].push_back(temp_v);
     g[temp_v].push_back(bcc_id+n);
     temp_v=par[temp_v];
    g[bcc_id+n].push_back(v);
    g[v].push_back(bcc_id+n);
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
void dfs(int u,int fa){
 if(u<=n){
```

for(int i=0;i<(int)g[u].size();i++){</pre>

```
int v=g[u][i];
   if(v==fa) continue;
   dfs(v,u);
   memset(tp,0x8f,sizeof tp);
   if(v<=n){
    tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
    tp[1]=max(
     dp[u][0]+dp[v][0]+1
     dp[u][1]+max(dp[v][0],dp[v][1])
   }else{
    tp[0]=dp[u][0]+dp[v][0];
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
   dp[u][0]=tp[0],dp[u][1]=tp[1];
 }else{
 for(int i=0;i<(int)g[u].size();i++){</pre>
  int v=g[u][i];
   if(v==fa) continue;
   dfs(v,u);
  min_dp[0][0]=0;
 min_dp[1][1]=1;
  min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
   if(v==fa) continue;
   memset(tmp,0x8f,sizeof tmp);
   tmp[0][0]=max(
    \min_{dp[0][0]+\max(dp[v][0],dp[v][1])}
    min_dp[0][1]+dp[v][0]
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
   tmp[1][0]=max(
    min_dp[1][0]+max(dp[v][0],dp[v][1]),
    min_dp[1][1]+dp[v][0]
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
   memcpy(min_dp,tmp,sizeof tmp);
  dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
  dp[u][0]=min_dp[0][0];
int main(){
int m,a,b;
scanf("%d%d",&n,&m);
for(int i=0;i<m;i++){</pre>
 scanf("%d%d",&a,&b);
 init_g[a].push_back(b);
  init_g[b].push_back(a);
par[1]=-1;
tarjan(1);
dfs(1,-1);
printf("%d\n", max(dp[1][0], dp[1][1]));
return 0;
8.7 Tree Knapsack
int dp[N][K]; PII obj[N];
vector<int> G[N];
void dfs(int u, int mx){
for(int s: G[u]) {
  if(mx < obj[s].first) continue;</pre>
  for(int i=0;i<=mx-obj[s].FF;i++)</pre>
   dp[s][i] = dp[u][i];
  dfs(s, mx - obj[s].first);
  for(int i=obj[s].FF;i<=mx;i++)</pre>
   dp[u][i] = max(dp[u][i],
    dp[s][i - obj[s].FF] + obj[s].SS);
}
}
8.8 N Queens Problem
vector< int > solve( int n ) {
 // no solution when n=2, 3
vector< int > ret;
if ( n % 6 == 2 ) {
  for ( int i = 2 ; i <= n ; i += 2 )</pre>
   ret.push_back( i );
```

```
ret.push_back( 3 ); ret.push_back( 1 );
  for (int i = 7 ; i \le n ; i += 2)
   ret.push_back( i );
  ret.push_back( 5 );
 } else if ( n % 6 == 3 ) {
  for ( int i = 4 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 2 );
  for ( int i = 5 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 1 ); ret.push_back( 3 );
 } else {
  for ( int i = 2 ; i <= n ; i += 2 )
   ret.push_back( i );
  for ( int i = 1 ; i <= n ; i += 2 )
   ret.push_back( i );
 return ret;
8.9 Binary Search On Fraction
struct Q {
 11 p, q;
 Q go(Q b, 11 d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 \le p,q \le N
Q frac_bs(ll N) {
 Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
 assert(pred(hi));
 bool dir = 1, L = 1, H = 1;
for (; L || H; dir = !dir) {
  11 len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)
if (Q mid = hi.go(lo, len + step);</pre>
     mid.p > N || mid.q > N || dir ^ pred(mid))
    t++:
   else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
 return dir ? hi : lo;
```