Contents			8 Misc 24 8.1 Theorems
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3	3.1 Euler Circuit	<b>5</b> 6	1 Basic
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4	4.1 Kuhn Munkres	10 10 11 11 11 12 12	<pre>//stack resize(change esp to rsp if 64-bit system) asm( "mov %0,%%esp\n" ::"g"(mem+10000000) ); // craziest way static void run_stack_sz(void(*func)(),size_t stsize){    char *stack, *send;    stack=(char *)malloc(stsize);    send=stack+stsize-16;</pre>
5	5.1 Prime Table       1         5.2 $\lfloor \frac{n}{i} \rfloor$ Enumeration       1         5.3 ax+by=gcd       1         5.4 Pollard Rho       1         5.5 Pi Count (Linear Sieve)       1         5.6 Range Sieve       1         5.7 Miller Rabin       1         5.8 Inverse Element       1         5.9 Euler Phi Function       1         5.10Gauss Elimination       1         5.11Fast Fourier Transform       1         5.12High Speed Linear Recurrence       1         5.13Chinese Remainder       1         5.14Berlekamp Massey       1         5.15NTT       1         5.16Polynomial Sqrt       1         5.17Polynomial Division       1	14 14 14 14 15 15 15 16 16 17 17	<pre>send=(char *)((uintptr_t)send/16*16); asm volatile(    "mov %%rsp, (%0)\n"    "mov %0, %%rsp\n"    :    : "r" (send)); func(); asm volatile(    "mov (%0), %%rsp\n"    :    : "r" (send)); free(stack); }  1.3 Pragma optimization  #pragma GCC optimize("Ofast, no-stack-protector") #pragma GCC optimize("no-math-errno, unroll-loops") #pragma GCC target("sse, sse2, sse3, ssse4") #pragma GCC target("popcnt, abm, mmx, avx, tune=native")</pre>
6	6.1 Point Class	18 18 18 19 19 19 20 20 20 21 21	<pre>#! /usr/bin/env python3 import subprocess as sp os_name =import('platform').system() cmd,prefix = [],"" if os_name == 'Windows':    cmd=["cmd", "/C"] else:    cmd = ["bash", "-c"]    prefix = "./"</pre>
7	7.1 Hash	21 21 22 22 23 23 23 23 23 23	<pre>def GetTestData(exe):    myout=sp.check_output(cmd+["%s%s"%(prefix, exe)])    return myout.decode("utf8") def Judge(a,b,testdata):    f = open("test.in", "w+")    f.write(testdata)    f.close()    c=sp.check_output(cmd+["%s%s<test.in"%(prefix, a)])="" b)])="" c="=" d='sp.check_output(cmd+["%s%s&lt;test.in"%(prefix,' d:<="" if="" not="" pre=""></test.in"%(prefix,></pre>

```
print("answer: %s"%c.decode("utf8"),end="")
print("output: %s"%d.decode("utf8"),end="")
print("WA!")
return False
return True
if __name__ == '__main__':
cnt = 0
isOK = True
while isOK:
cnt += 1
print(cnt)
isOK=Judge("sol", "mysol", GetTestData("gen"))
```

## 1.5 Quick Random

```
template < class T,T x1,T x2,T x3,int y1,int y2,int y3>
struct PRNG {
  using S = typename std::make_signed<T>::type;
  PRNG(T _s = 0) : s(_s) {}
  T next() {
   T z = (s += x1);
    z = (z ^ (z >> y1)) * x2;
    z = (z ^ (z >> y2)) * x3;
    return z ^ (z >> y3);
 T next(T n) { return next() % n; }
 S next(S 1, S r){return l+next(r-l+1);}
 T operator()() { return next(); }
 T operator()(T n) { return next(n); }
 S operator()(S 1, S r) { return next(1, r); }
  static T gen(T s) { return PRNG(s)(); }
  template < class U>
  void shuffle(U first,U last){
    size_t n=last-first;
    for(size_t i=0;i<n;i++)</pre>
      swap(first[i],first[next(i+1)]);
 }
};
using R32=PRNG<uint32_t,0x9E3779B1,0x85EBCA6B,</pre>
0xC2B2AE35,16,13,16>;
R32 r32:
using R64=PRNG<uint64_t,0x9E3779B97F4A7C15,</pre>
0xBF58476D1CE4E5B9,0x94D049BB133111EB,30,27,31>;
```

# 1.6 IO Optimization

```
static inline int gc() {
  static char buf[ 1 << 20 ], *p = buf, *end = buf;</pre>
  if ( p == end ) {
    end = buf + fread( buf, 1, 1 << 20, stdin );
    if ( end == buf ) return EOF;
    p = buf;
  return *p++;
template < typename T >
static inline bool gn( T &_ ) {
  register int c = gc(); register T __ = 1; _ = 0;
  while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
if(c == '-') { __ = -1; c = gc(); }
  if(c == EOF) return false;
  while('0'<=c&c<='9') _{-} = _{-} * 10 + c - '0', c = gc();
    *= __;
  return true;
template < typename T, typename ...Args >
static inline bool gn( T &x, Args &...args )
{ return gn(x) && gn(args...); }
```

#### 2 Data Structure

#### 2.1 Bigint

```
class BigInt{
  private:
    using lld = int_fast64_t;
    #define PRINTF_ARG PRIdFAST64
    #define LOG_BASE_STR "9"
```

```
static constexpr lld BASE = 1000000000;
 static constexpr int LOG_BASE = 9;
 vector<lld> dig;
 bool neg;
 inline int len() const { return (int) dig.size(); }
 inline int cmp_minus(const BigInt& a) const {
   if(len() == 0 && a.len() == 0) return 0;
if(neg ^ a.neg)return (int)a.neg*2 - 1;
   if(len()!=a.len())
     return neg?a.len()-len():len()-a.len();
   for(int i=len()-1;i>=0;i--) if(dig[i]!=a.dig[i])
     return neg?a.dig[i]-dig[i]:dig[i]-a.dig[i];
   return 0;
 inline void trim(){
   while(!dig.empty()&&!dig.back())dig.pop_back();
   if(dig.empty()) neg = false;
public:
 BigInt(): dig(vector<lld>()), neg(false){}
 BigInt(lld a): dig(vector<lld>()){
   neg = a<0; dig.push_back(abs(a));</pre>
   trim();
 BigInt(const string& a): dig(vector<lld>()){
   assert(!a.empty()); neg = (a[0]=='-');
   for(int i=((int)a.size())-1;i>=neg;i-=LOG_BASE){
     11d cur = 0;
     for(int j=min(LOG_BASE-1,i-neg);j>=0;j--)
       cur = cur*10+a[i-j]-'0';
     dig.push_back(cur);
   } trim();
 inline bool operator<(const BigInt& a)const</pre>
   {return cmp_minus(a)<0;}
 inline bool operator<=(const BigInt& a)const</pre>
   {return cmp_minus(a)<=0;}
 inline bool operator==(const BigInt& a)const
   {return cmp_minus(a)==0;}
 inline bool operator!=(const BigInt& a)const
   {return cmp_minus(a)!=0;}
 inline bool operator>(const BigInt& a)const
   {return cmp_minus(a)>0;}
 inline bool operator>=(const BigInt& a)const
   {return cmp_minus(a)>=0;}
 BigInt operator-() const {
   BigInt ret = *this;
   ret.neg ^= 1;
   return ret;
 BigInt operator+(const BigInt& a) const {
   if(neg) return -(-(*this)+(-a));
   if(a.neg) return (*this)-(-a);
   int n = max(a.len(), len());
   BigInt ret; ret.dig.resize(n);
   11d pro = 0;
   for(int i=0;i<n;i++) {</pre>
     ret.dig[i] = pro;
     if(i < a.len()) ret.dig[i] += a.dig[i];</pre>
     if(i < len()) ret.dig[i] += dig[i];</pre>
     pro = 0;
     if(ret.dig[i] >= BASE) pro = ret.dig[i]/BASE;
     ret.dig[i] -= BASE*pro;
   if(pro != 0) ret.dig.push_back(pro);
   return ret;
 BigInt operator-(const BigInt& a) const {
   if(neg) return -(-(*this) - (-a));
   if(a.neg) return (*this) + (-a);
   int diff = cmp_minus(a);
   if(diff < 0) return -(a - (*this));</pre>
   if(diff == 0) return 0;
   BigInt ret; ret.dig.resize(len(), 0);
   for(int i=0;i<len();i++) {</pre>
     ret.dig[i] += dig[i];
     if(i < a.len()) ret.dig[i] -= a.dig[i];</pre>
     if(ret.dig[i] < 0){
       ret.dig[i] += BASE;
       ret.dig[i+1]--;
```

```
ret.trim():
  return ret;
BigInt operator*(const BigInt& a) const {
  if(!len()||!a.len()) return 0;
  BigInt ret; ret.dig.resize(len()+a.len()+1);
  ret.neg = neg ^ a.neg;
  for(int i=0;i<len();i++)</pre>
    for(int j=0;j<a.len();j++){
  ret.dig[i+j] += dig[i] * a.dig[j];</pre>
      if(ret.dig[i+j] >= BASE) {
        lld x = ret.dig[i+j] / BASE;
        ret.dig[i+j+1] += x;
        ret.dig[i+j] -= x * BASE;
      }
    }
  ret.trim():
  return ret;
BigInt operator/(const BigInt& a) const {
  assert(a.len());
  if(len() < a.len()) return 0;</pre>
  BigInt ret; ret.dig.resize(len()-a.len()+1);
  ret.neg = a.neg;
  for(int i=len()-a.len();i>=0;i--){
    11d l = 0, r = BASE;
    while (r-l > 1) {
      lld mid = (1+r)>>1;
      ret.dig[i] = mid;
      if(ret*a<=(neg?-(*this):(*this))) l = mid;</pre>
      else r = mid:
    ret.dig[i] = 1;
  ret.neg ^= neg; ret.trim();
  return ret;
BigInt operator%(const BigInt& a) const {
  return (*this) - (*this) / a * a;
friend BigInt abs(BigInt a){
  a.neg = 1; return a;
friend void swap(BigInt& a, BigInt& b){
  swap(a.dig, b.dig); swap(a.neg, b.neg);
friend istream& operator>>(istream& ss, BigInt& a){
  string s; ss >> s; a = s;
  return ss:
friend ostream&operator<<(ostream&o, const BigInt&a){</pre>
  if(a.len() == 0) return o << '0';</pre>
  if(a.neg) o <<</pre>
  ss << o.dig.back();</pre>
  for(int i=a.len()-2;i>=0;i--)
    o<<setw(LOG_BASE)<<setfill('0')<<a.dig[i];</pre>
  return o;
inline void print() const {
  if(len() == 0){putchar('0'); return;}
  if(neg) putchar('-');
  printf("%" PRINTF_ARG, dig.back());
  for(int i=len()-2;i>=0;i--)
    printf("%0" LOG_BASE_STR PRINTF_ARG, dig[i]);
#undef PRINTF_ARG
#undef LOG_BASE_STR
```

#### 2.2 Dark Magic

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using __gnu_pbds::pairing_heap_tag;
using __gnu_pbds::binary_heap_tag;
using __gnu_pbds::binomial_heap_tag;
using __gnu_pbds::rc_binomial_heap_tag;
using __gnu_pbds::thin_heap_tag;
using __gnu_pbds::thin_heap_tag;
template<typename T>
using pbds_heap=__gnu_pbds::prioity_queue<T,less<T>,\
pairing_heap_tag>;
```

```
using __gnu_pbds::rb_tree_tag;
using __gnu_pbds::ov_tree_tag;
using __gnu_pbds::splay_tree_tag;
template<typename T>
using ordered_set = __gnu_pbds::tree<T,\</pre>
__gnu_pbds::null_type,less<T>,rb_tree_tag,\
__gnu_pbds::tree_order_statistics_node_update>;
template<typename A, typename B>
using hTable1=__gnu_pbds::cc_hash_table<A,B>;
template<typename A, typename B>
using hTable2=__gnu_pbds::gp_hash_table<A,B>;
int main(){
  ordered_set<int> ss;
  ss.insert(1); ss.insert(5);
  assert(*ss.find_by_order(0)==1);
  assert(ss.order_of_key(-1)==0);
  pbds_heap pq1, pq2;
  pq1.push(1); pq2.push(2);
  pq1.join(pq2);
  assert(pq2.size()==0);
  auto it = pq1.push(87);
  pq1.modify(it, 19);
  return 0;
```

### 2.3 SkewHeap

```
template < typename T, typename cmp = less< T > >
class SkewHeap{
private:
  struct SkewNode{
    T x;
    SkewNode *lc, *rc;
    SkewNode( T a = 0 ) : x(a), lc(0), rc(0) {}
  } *root;
  cmp CMP_;
  size t count;
  SkewNode* Merge( SkewNode* a, SkewNode* b ) {
    if ( !a or !b ) return a ? a : b;
    if ( CMP_( a->x, b->x ) ) swap( a, b );
    a -> rc = Merge( a->rc, b );
    swap( a -> lc, a->rc );
    return a;
public:
  SkewHeap(): root( 0 ), count( 0 ) {}
  size_t size() { return count; }
  bool empty() { return count == 0; }
  T top() { return root->x; }
  void clear(){ root = 0; count = 0; }
  void push ( const T& x ) {
    SkewNode* a = new SkewNode( x );
    count += 1; root = Merge( root, a );
  void join( SkewHeap& a ) {
    count += a.count; a.count = 0;
    root = Merge( root, a.root );
  void pop() {
    count--; root = Merge( root->lc, root->rc );
  friend void swap( SkewHeap& a, SkewHeap& b ) {
    swap( a.root, b.root ); swap( a.count, b.count );
};
```

# 2.4 Disjoint Set

```
class DJS{
private:
    vector< int > fa, sz, sv;
    vector< pair< int*, int > > opt;
    inline void assign( int *k, int v ) {
        opt.emplace_back( k, *k );
        *k = v;
    }
public:
    inline void init( int n ) {
        fa.resize( n ); iota( fa.begin(), fa.end(), 0 );
        sz.resize( n ); fill( sz.begin(), sz.end(), 1 );
        opt.clear();
```

```
int query( int x ) {
    return ( fa[ x ] == x ) ? x : query( fa[ x ] );
  inline void merge( int a, int b ) {
    int af = query( a ), bf = query( b );
    if( af == bf ) return;
    if( sz[ af ] < sz[ bf ] ) swap( af, bf );</pre>
    assign( &fa[ bf ], fa[ af ] );
    assign( &sz[ af ], sz[ af ] + sz[ bf ] );
  inline void save() {sv.push_back( (int)opt.size() );}
  inline void undo() {
    int ls = sv.back(); sv.pop_back();
    while ( ( int ) opt.size() > ls ) {
      pair< int*, int > cur = opt.back();
      *cur.first = cur.second;
      opt.pop_back();
  }
};
```

#### 2.5 Link-Cut Tree

```
struct Node{
 Node *par, *ch[2];
 int xor_sum,v;
  bool is_rev;
 Node(int _v){
   v=xor_sum=_v;is_rev=false;
    par=ch[0]=ch[1]=nullptr;
 inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
  inline void down(){
   if(is_rev){
      if(ch[0]!=nullptr) ch[0]->set_rev();
      if(ch[1]!=nullptr) ch[1]->set_rev();
      is_rev=false;
   }
 inline void up(){
   xor_sum=v;
    if(ch[0]!=nullptr){
      xor_sum^=ch[0]->xor_sum;
      ch[0]->par=this;
    if(ch[1]!=nullptr){
      xor_sum^=ch[1]->xor_sum;
      ch[1]->par=this;
   }
  inline bool is_root(){
    return par==nullptr ||\
      (par->ch[0]!=this && par->ch[1]!=this);
  bool is_rch(){return !is_root() && par->ch[1]==this;}
} *node[maxn],*stk[maxn];
int top;
void to_child(Node* p,Node* c,bool dir){
 p->ch[dir]=c;
 p->up();
inline void rotate(Node* node){
 Node* par=node->par;
Node* par_par=par->par;
  bool dir=node->is_rch();
 bool par_dir=par->is_rch();
 to_child(par,node->ch[!dir],dir);
  to_child(node,par,!dir);
 if(par_par!=nullptr && par_par->ch[par_dir]==par)
    to_child(par_par,node,par_dir);
  else node->par=par_par;
inline void splay(Node* node){
 Node* tmp=node;
  stk[top++]=node;
 while(!tmp->is_root()){
   tmp=tmp->par;
    stk[top++]=tmp;
 while(top) stk[--top]->down();
```

```
for(Node *fa=node->par;
   !node->is_root();
   rotate(node), fa=node->par)
    if(!fa->is_root())
      rotate(fa->is_rch()==node->is_rch()?fa:node);
inline void access(Node* node){
  Node* last=nullptr;
  while(node!=nullptr){
    splay(node);
    to_child(node, last, true);
    last=node;
    node=node->par;
  }
inline void change_root(Node* node){
  access(node);splay(node);node->set_rev();
inline void link(Node* x,Node* y){
  change_root(x);splay(x);x->par=y;
inline void split(Node* x,Node* y){
  change_root(x);access(y);splay(x);
  to_child(x,nullptr,true);y->par=nullptr;
inline void change_val(Node* node,int v){
  access(node);splay(node);node->v=v;node->up();
inline int query(Node* x,Node* y){
  change_root(x);access(y);splay(y);
  return y->xor_sum;
inline Node* find_root(Node* node){
  access(node);splay(node);
  Node* last=nullptr;
  while(node!=nullptr){
    node->down();last=node;node=node->ch[0];
  return last;
set<pii> dic;
inline void add_edge(int u,int v){
  if(u>v) swap(u,v);
  if(find_root(node[u])==find_root(node[v])) return;
  dic.insert(pii(u,v));
  link(node[u],node[v]);
inline void del_edge(int u,int v){
  if(u>v) swap(u,v);
  if(dic.find(pii(u,v))==dic.end()) return;
  dic.erase(pii(u,v));
  split(node[u],node[v]);
```

#### 2.6 LiChao Segment Tree

```
struct Line{
  int m, k, id;
  Line() : id( -1 ) {}
  Line(int a, int b, int c)
: m(a), k(b), id(c) {}
int at(int x) { return m * x + k; }
};
class LiChao {
  private:
    int n; vector< Line > nodes;
    inline int lc( int x ) { return 2 * x + 1; }
inline int rc( int x ) { return 2 * x + 2; }
     void insert( int 1, int r, int id, Line ln ) {
       int m = (1 + r) >> 1;
       if ( nodes[ id ].id == -1 ) {
         nodes[ id ] = ln;
         return;
       bool atLeft = nodes[ id ].at( l ) < ln.at( l );</pre>
       if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
         atLeft ^= 1; swap( nodes[ id ], ln );
       if ( r - l == 1 ) return;
       if ( atLeft ) insert( l, m, lc( id ), ln );
       else insert( m, r, rc( id ), ln );
```

```
int query( int 1, int r, int id, int x ) {
      int ret = 0;
      if ( nodes[ id ].id != -1 )
        ret = nodes[ id ].at( x );
      int m = ( 1 + r ) >> 1;
      if ( r - l == 1 ) return ret;
      else if ( x < m )</pre>
        return max( ret, query( 1, m, lc( id ), x ) );
      else
        return max( ret, query( m, r, rc( id ), x ) );
  public:
    void build( int n_ ) {
      n = n_; nodes.clear();
      nodes.resize( n << 2, Line() );</pre>
    void insert( Line ln ) { insert( 0, n, 0, ln ); }
    int query( int x ) { return query( 0, n, 0, x ); }
} lichao;
```

# 2.7 Treap

```
namespace Treap{
  #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
  struct node{
    int size;
    uint32_t pri;
    node *lc, *rc;
    node() : size(0), pri(rand()), lc( 0 ), rc( 0 ) {}
    void pull() {
      size = 1;
      if ( lc ) size += lc->size;
      if ( rc ) size += rc->size;
    }
  };
  node* merge( node* L, node* R ) {
    if ( not L or not R ) return L ? L : R;
    if ( L->pri > R->pri ) {
      L->rc = merge( L->rc, R ); L->pull();
      return L;
    } else {
      R->lc = merge( L, R->lc ); R->pull();
      return R;
    }
  void split_by_size( node*rt,int k,node*&L,node*&R ) {
    if ( not rt ) L = R = nullptr;
    else if( sz( rt->lc ) + 1 <= k ) {
     L = rt;
      split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
      L->pull();
    } else {
      R = rt;
      split_by_size( rt->lc, k, L, R->lc );
      R->pull();
    }
  #undef sz
}
```

# 2.8 SparseTable

```
template < typename T, typename Cmp_ = less< T > >
class SparseTable {
private:
 vector< vector< T > > tbl;
  vector< int > lg;
 T cv( T a, T b ) {
    return Cmp_()( a, b ) ? a : b;
public:
 void init( T arr[], int n ) {
    // 0-base
    lg.resize(n+1);
    lg[0] = -1;
    for( int i=1 ; i<=n ; ++i ) lg[i] = lg[i>>1] + 1;
    tbl.resize( lg[n] + 1 );
    tbl[ 0 ].resize( n );
    copy( arr, arr + n, tbl[ 0 ].begin() );
    for ( int i = 1 ; i <= lg[ n ] ; ++ i ) {</pre>
```

```
int len = 1 << ( i - 1 ), sz = 1 << i;
tbl[ i ].resize( n - sz + 1 );
for ( int j = 0 ; j <= n - sz ; ++ j )
    tbl[i][j] = cv(tbl[i-1][j], tbl[i-1][j+len]);
}

T query( int l, int r ) {
    // 0-base [L, r)
    int wh = lg[ r - l ], len = 1 << wh;
    return cv( tbl[ wh ][ l ], tbl[ wh ][ r - len ] );
}
};</pre>
```

#### 2.9 Linear Basis

```
struct LinearBasis {
private:
  int n, sz;
   vector< llu > B;
   inline llu two( int x ){ return ( ( llu ) 1 ) << x; }</pre>
public:
   void init( int n_ ) {
    n = n_; B.clear(); B.resize( n ); sz = 0;
   void insert( llu x ) {
     // add x into B
     for ( int i = n-1; i >= 0 ; --i ) if( two(i) & x ){
       if ( B[ i ] ) x ^= B[ i ];
       else
         B[i] = x; sz++;
         for ( int j = i - 1 ; j >= 0 ; -- j )
            if( B[ j ] && ( two( j ) & B[ i ] ) )
B[ i ] ^= B[ j ];
         for (int j = i + 1; j < n; ++ j)
  if ( two( i ) & B[ j ] )
   B[ j ] ^= B[ i ];</pre>
         break;
       }
     }
   inline int size() { return sz; }
   bool check( llu x ) {
     // is x in span(B) ?
     for ( int i = n-1; i >= 0; --i ) if( two(i) & x )
       if( B[ i ] ) x ^= B[ i ];
       else return false;
     return true;
   llu kth_small(llu k) {
     /** 1-base would always > 0 **/
     /** should check it **/
     /* if we choose at least one element
        but size(B)(vectors in B)==N(original elements)
        then we can't get 0 */
     llu ret = 0;
     for ( int i = 0 ; i < n ; ++ i ) if( B[ i ] ) {</pre>
       if( k & 1 ) ret ^= B[ i ];
       k >>= 1;
     return ret;
} base;
```

# 3 Graph

#### 3.1 Euler Circuit

```
bool vis[ N ]; size_t la[ K ];
void dfs( int u, vector< int >& vec ) {
  while ( la[ u ] < G[ u ].size() ) {
    if( vis[ G[ u ][ la[ u ] ].second ] ) {
        ++ la[ u ];
        continue;
    }
    int v = G[ u ][ la[ u ] ].first;
    vis[ G[ u ][ la[ u ] ].second ] = true;
    ++ la[ u ]; dfs( v, vec );
    vec.push_back( v );
  }
}</pre>
```

# 3.2 BCC Edge

```
class BCC{
private:
  vector< int > low, dfn;
  int cnt;
  vector< bool > bridge;
  vector< vector< PII > > G;
  void dfs( int w, int f ) {
    low[ w ] = dfn[ w ] = cnt++;
    for ( auto [ u, t ] : G[ w ] ) {
  if ( u == f ) continue;
      if ( dfn[ u ] != 0 ) {
        low[ w ] = min( low[ w ], dfn[ u ] );
      }else{
        dfs( u, w );
        low[w] = min(low[w], low[u]);
        if ( low[ u ] > dfn[ w ] ) bridge[ t ] = true;
   }
public:
 void init( int n, int m ) {
    G.resize(n); cnt = 0;
    fill( G.begin(), G.end(), vector< PII >() );
    bridge.clear(); bridge.resize( m );
    low.clear(); low.resize( n );
    dfn.clear(); dfn.resize( n );
  void add_edge( int u, int v ) {
    // should check for multiple edge
    G[ u ].emplace_back( v, cnt );
   G[ v ].emplace_back( u, cnt ++ );
  void solve(){ cnt = 1; dfs( 0, 0 ); }
  // the id will be same as insert order, 0-base
  bool is_bridge( int x ) { return bridge[ x ]; }
```

#### 3.3 BCC Vertex

```
class BCC{
  private:
    int n, ecnt;
    vector< vector< pair< int, int > > > G;
    vector< int > low, dfn, id;
    vector< bool > vis, ap;
    void dfs( int u, int f, int d ) {
      int child = 0;
      dfn[ u ] = low[ u ] = d; vis[ u ] = true;
      for ( auto e : G[ u ] ) if ( e.first != f ) {
        if ( vis[ e.first ] ) {
          low[ u ] = min( low[ u ], dfn[ e.first ] );
        } else {
          dfs( e.first, u, d + 1 ); child ++;
low[ u ] = min( low[ u ], low[ e.first ] );
          if ( low[ e.first ] >= d ) ap[ u ] = true;
        }
      if ( u == f and child <= 1 ) ap[ u ] = false;</pre>
    void mark( int u, int idd ) {
      // really??????????
      if ( ap[ u ] ) return;
      for ( auto e : G[ u ] )
        if( id[ e.second ] != -1 ) {
          id[ e.second ] = idd;
          mark( e.first, idd );
  public:
    void init( int n_ ) {
      ecnt = 0, n = n_{j};
      G.clear(); G.resize( n );
      low.resize( n ); dfn.resize( n );
      ap.clear(); ap.resize( n );
      vis.clear(); vis.resize( n );
    void add_edge( int u, int v ) {
      G[ u ].emplace_back( v, ecnt );
      G[ v ].emplace_back( u, ecnt ++ );
```

```
}
void solve() {
    for ( int i = 0 ; i < n ; ++ i )
        if ( not vis[ i ] ) dfs( i, i, 0 );
    id.resize( ecnt );
    fill( id.begin(), id.end(), -1 );
    ecnt = 0;
    for ( int i = 0 ; i < n ; ++ i )
        if ( ap[ i ] ) for ( auto e : G[ i ] )
        if( id[ e.second ] != -1 ) {
            id[ e.second ] = ecnt;
            mark( e.first, ecnt ++ );
        }
}
int get_id( int x ) { return id[ x ]; }
int count() { return ecnt; }
bool is_ap( int u ) { return ap[ u ]; }
}
bcc;</pre>
```

# 3.4 2-SAT (SCC)

```
class TwoSat{
  private:
    int n;
    vector<vector<int>> rG,G,sccs;
    vector<int> ord,idx;
    vector<bool> vis,result;
    void dfs(int u){
      vis[u]=true;
      for(int v:G[u])
        if(!vis[v]) dfs(v);
      ord.push_back(u);
    void rdfs(int u){
      vis[u]=false;idx[u]=sccs.size()-1;
      sccs.back().push_back(u);
      for(int v:rG[u])
        if(vis[v])rdfs(v);
  public:
    void init(int n_){
      n=n_;G.clear();G.resize(n);
      rG.clear();rG.resize(n);
      sccs.clear();ord.clear();
      idx.resize(n);result.resize(n);
    void add_edge(int u,int v){
      G[u].push_back(v);rG[v].push_back(u);
    void orr(int x,int y){
      if ((x^y)==1)return;
      add_edge(x^1,y); add_edge(y^1,x);
    bool solve(){
      vis.clear();vis.resize(n);
      for(int i=0;i<n;++i)</pre>
        if(not vis[i])dfs(i);
      reverse(ord.begin(),ord.end());
      for (int u:ord){
        if(!vis[u])continue;
        sccs.push_back(vector<int>());
        rdfs(u);
      for(int i=0;i<n;i+=2)</pre>
        if(idx[i]==idx[i+1])
          return false;
      vector<bool> c(sccs.size());
      for(size_t i=0;i<sccs.size();++i){</pre>
        for(size_t j=0;j<sccs[i].size();++j){</pre>
          result[sccs[i][j]]=c[i];
          c[idx[sccs[i][j]^1]]=!c[i];
        }
      }
      return true;
    bool get(int x){return result[x];}
    inline int get_id(int x){return idx[x];}
    inline int count(){return sccs.size();}
} sat2;
```

### 3.5 Lowbit Decomposition

```
class LowbitDecomp{
private:
 int time_, chain_, LOG_N;
  vector< vector< int > > G, fa;
  vector< int > tl, tr, chain, chain_st;
 // chain_ : number of chain
 // tl, \overline{\mathsf{tr}}[\mathsf{u}] : subtree interval in the seq. of \mathsf{u}
 // chain_st[ u ] : head of the chain contains u
 // chian[ u ] : chain id of the chain u is on
 inline int lowbit( int x ) {
    return x & ( -x );
  void predfs( int u, int f ) {
    chain[ u ] = 0;
    for ( int v : G[ u ] ) {
   if ( v == f ) continue;
      predfs( v, u );
      if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
        chain[ u ] = chain[ v ];
    if ( not chain[ u ] )
      chain[ u ] = chain_ ++;
  void dfschain( int u, int f ) {
    fa[ u ][ 0 ] = f;
    for ( int i = 1 ; i < LOG_N ; ++ i )</pre>
      fa[ u ][ i ] = fa[ fa[ u ][ i - 1 ] ][ i - 1 ];
    tl[ u ] = time_++;
    if ( not chain_st[ chain[ u ] ] )
      chain_st[ chain[ u ] ] = u;
    for ( int v : G[ u ] )
  if ( v != f and chain[ v ] == chain[ u ] )
        dfschain( v, u );
    for ( int v : G[ u ] )
      if ( v != f and chain[ v ] != chain[ u ] )
        dfschain( v, u );
    tr[ u ] = time_;
  inline bool anc( int u, int v ) {
    return tl[ u ] <= tl[ v ] \</pre>
      and tr[ v ] <= tr[ u ];</pre>
public:
 inline int lca( int u, int v ) {
    if ( anc( u, v ) ) return u;
    for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
      if ( not anc( fa[ u ][ i ], v ) )
        u = fa[ u ][ i ];
    return fa[ u ][ 0 ];
  void init( int n ) {
    for ( LOG_N = 0 ; ( 1 << LOG_N ) < n ; ++ LOG_N );
    fa.clear();
    fa.resize( n, vector< int >( LOG_N ) );
    G.clear(); G.resize( n );
    tl.clear(); tl.resize( n );
tr.clear(); tr.resize( n );
    chain.clear(); chain.resize( n );
    chain_st.clear(); chain_st.resize( n );
  void add_edge( int u , int v ) {
    // 1-base
    G[ u ].push_back( v );
    G[ v ].push_back( u );
  void decompose(){
    chain_ = 1;
    predfs( 1, 1 );
    time_{-} = 0;
    dfschain( 1, 1 );
  PII get_inter( int u ) { return {tl[ u ], tr[ u ]}; }
  vector< PII > get_path( int u , int v ){
    vector< PII > res;
    int g = lca( u, v );
while ( chain[ u ] != chain[ g ] ) {
      int s = chain_st[ chain[ u ] ];
      res.emplace_back( tl[ s ], tl[ u ] + 1 );
      u = fa[ s ][ 0 ];
```

```
}
res.emplace_back( tl[ g ], tl[ u ] + 1 );
while ( chain[ v ] != chain[ g ] ) {
    int s = chain_st[ chain[ v ] ];
    res.emplace_back( tl[ s ], tl[ v ] + 1 );
    v = fa[ s ][ 0 ];
}
res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
return res;
/* res : list of intervals from u to v
    * ( note only nodes work, not edge )
    * usage :
    * vector< PII >& path = tree.get_path( u , v )
    * for( auto [ l, r ] : path ) {
        * 0-base [ l, r )
        * }
        */
}
tree;
```

# 3.6 MaxClique

```
// contain a self loop u to u, than u won't in clique
template < size t MAXN >
class MaxClique{
  private:
    using bits = bitset< MAXN >;
    bits popped, G[ MAXN ], ans;
    size_t deg[ MAXN ], deo[ MAXN ], n;
    inline void sort_by_degree() {
      popped.reset();
      for ( size_t i = 0 ; i < n ; ++ i )
          deg[ i ] = G[ i ].count();
      for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
           size_t mi = MAXN, id = 0;
           for ( size_t j = 0 ; j < n ; ++ j )
    if ( not popped[ j ] and deg[ j ] < mi )</pre>
                   mi = deg[id = j];
           popped[ deo[ i ] = id ] = 1;
           for( size_t u = G[ i ]._Find_first() ;
            u < n ; u = G[ i ]._Find_next( u ) )
               -- deg[ u ];
    void BK( bits R, bits P, bits X ) {
      if ( R.count() + P.count() <= ans.count() )return</pre>
      if ( not P.count() and not X.count() ) {
        if ( R.count() > ans.count() ) ans = R;
      }
      /* greedily chosse max degree as pivot
      bits cur = P | X; size_t pivot = 0, sz = 0;
      for ( size_t u = cur._Find_first() ;
       u < n ; u = cur._Find_next( u ) )
    if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
      cur = P & ( \sim G[ pivot ] );
      */ // or simply choose first
      bits cur = P & ( \sim G[ (P \mid X)._Find_first() ] );
      for ( size_t u = cur._Find_first() ;
       u < n ; u = cur._Find_next( u ) ) {
        if ( R[ u ] ) continue;
        R[u] = 1;
        BK( R, P & G[ u ], X & G[ u ]);
        R[u] = P[u] = 0, X[u] = 1;
      }
  public:
    void init( size_t n_ ) {
      n = n_;
      for ( size_t i = 0 ; i < n ; ++ i )</pre>
        G[ i ].reset();
      ans.reset();
    void add_edges( int u, bits S ) { G[ u ] = S; }
    void add_edge( int u, int v ) {
      G[u][v] = G[v][u] = 1;
    int solve() {
      sort_by_degree(); // or simply iota( deo... )
      for ( size_t i = 0 ; i < n ; ++ i )</pre>
```

```
deg[ i ] = G[ i ].count();
       bits pob, nob = 0; pob.set();
       for ( size_t i = n ; i < MAXN ; ++ i ) pob[ i</pre>
           1=0:
       for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
         size_t v = deo[ i ];
         bits tmp; tmp[ v ] = 1;
         BK( tmp, pob & G[ v ], nob & G[ v ]);
         pob[ v ] = 0, nob[ v ] = 1;
       return static_cast< int >( ans.count() );
     }
|};
```

# 3.7 Virtural Tree

```
inline bool cmp(const int &i, const int &j) {
 return dfn[i] < dfn[j];</pre>
void build(int vectrices[], int k) {
  static int stk[MAX_N];
  sort(vectrices, vectrices + k, cmp);
  stk[sz++] = 0;
  for (int i = 0; i < k; ++i) {</pre>
    int u = vectrices[i], lca = LCA(u, stk[sz - 1]);
    if (lca == stk[sz - 1]) stk[sz++] = u;
    else {
      while (sz \ge 2 \&\& dep[stk[sz - 2]] \ge dep[lca]) {
        addEdge(stk[sz - 2], stk[sz - 1]);
        SZ - -:
      if (stk[sz - 1] != lca) {
        addEdge(lca, stk[--sz]);
        stk[sz++] = lca, vectrices[cnt++] = lca;
      stk[sz++] = u;
   }
  for (int i = 0; i < sz - 1; ++i)</pre>
    addEdge(stk[i], stk[i + 1]);
```

#### 3.8 Tree Hashing

```
uint64_t hsah( int u, int f ) {
    uint64_t r = 127;
    for ( int v : G[ u ] ) {
        if ( v == f ) continue;
        uint64_t hh = hsah( v, u );
        r = r + (hh * hh) \% mod;
    return r;
}
```

# 3.9 Minimum Mean Cycle

```
/* minimum mean cycle O(VE) */
struct MMC{
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
#define V 1021
#define inf 1e9
 struct Edge { int v,u; double c; };
  int n, m, prv[V][V], prve[V][V], vst[V];
 Edge e[E];
 vector<int> edgeID, cycle, rho;
 double d[V][V];
 void init( int _n ) { n = _n; m = 0; }
 // WARNING: TYPE matters
 void add_edge( int vi , int ui , double ci )
  { e[ m ++ ] = { vi , ui , ci }; }
  void bellman_ford() {
    for(int i=0; i<n; i++) d[0][i]=0;</pre>
    for(int i=0; i<n; i++) {</pre>
      fill(d[i+1], d[i+1]+n, inf);
      for(int j=0; j<m; j++) {</pre>
        int v = e[j].v, u = e[j].u;
        if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
          d[i+1][u] = d[i][v]+e[j].c;
          prv[i+1][u] = v;
```

```
prve[i+1][u] = j;
        }
      }
    }
  double solve(){
    // returns inf if no cycle, mmc otherwise
    double mmc=inf;
    int st = -1:
    bellman_ford();
    for(int i=0; i<n; i++) {</pre>
      double avg=-inf;
      for(int k=0; k<n; k++) {</pre>
        if(d[n][i]<inf-eps)</pre>
          avg=max(avg,(d[n][i]-d[k][i])/(n-k));
        else avg=max(avg,inf);
      if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
    FZ(vst); edgeID.clear(); cycle.clear(); rho.clear()
    for (int i=n; !vst[st]; st=prv[i--][st]) {
      vst[st]++;
      edgeID.PB(prve[i][st]);
      rho.PB(st);
    while (vst[st] != 2) {
      int v = rho.back(); rho.pop_back();
      cycle.PB(v);
      vst[v]++;
    reverse(ALL(edgeID));
    edgeID.resize(SZ(cycle));
    return mmc;
} mmc;
```

# 3.10 Mo's Algorithm on Tree

```
int n, q, nxt[ N ], to[ N ], hd[ N ];
struct Que{
  int u, v, id;
} que[ N ];
void init() {
  cin >> n >> q;
  for ( int i = 1 ; i < n ; ++ i ) {</pre>
    int u, v; cin >> u >> v;
    nxt[ i << 1 | 0 ] = hd[ u ];</pre>
    to[ i << 1 | 0 ] = v;
    hd[u] = i << 1 | 0;
    nxt[ i << 1 | 1 ] = hd[ v ];</pre>
    to[ i << 1 | 1 ] = u;
    hd[ v ] = i << 1 | 1;
  for ( int i = 0 ; i < q ; ++ i ) {</pre>
    cin >> que[ i ].u >> que[ i ].v; que[ i ].id = i;
int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
void dfs( int u, int f ) {
  dfn[ u ] = dfn_++; int saved_rbp = stk_;
for ( int v_ = hd[ u ] ; v_ ; v_ = nxt[ v_ ] ) {
    if ( to[ v_ ] == f ) continue;
    dfs( to[ v_ ], u );
    if ( stk_ - saved_rbp < SQRT_N ) continue;</pre>
    for ( ++ block_ ; stk_ != saved_rbp ; )
       block_id[ stk[ -- stk_ ] ] = block_;
  stk[ stk_ ++ ] = u;
bool inPath[ N ];
void Diff( int u ) {
  if ( inPath[ u ] ^= 1 )
    // remove this edge
  else
    // add this edge
void traverse( int& origin_u, int u ) {
  for ( int g = lca( origin_u, u );
    origin_u != g ; origin_u = parent_of[ origin_u ] )
```

Diff( origin\_u );

```
( int v = u ; v != origin_u ; v = parent_of[ v ]
   Diff( v );
 origin_u = u;
void solve() {
 dfs( 1, 1 );
  while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
  sort( que, que + q, [] ( const Que& x, const Que& y )
    return tie( block_id[ x.u ], dfn[ x.v ] )
            < tie( block_id[ y.u ], dfn[ y.v ] );</pre>
 } );
  int U = 1, V = 1;
 for ( int i = 0 ; i < q ; ++ i ) {</pre>
   pass( U, que[ i ].u );
    pass( V, que[ i ].v );
    // we could get our answer of que[ i ].id
}
Method 2:
dfs u:
 push u
 iterate subtree
 push u
Let P = LCA(u, v), and St(u) <= St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
```

#### 3.11 Minimum Steiner Tree

```
// Minimum Steiner Tree
// O(V 3^T + V^2 2^T)
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
  int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
  void init( int _n ){
    n = _n;
for( int i = 0 ; i < n ; i ++ ){</pre>
       for( int j = 0 ; j < n ; j ++ )</pre>
       dst[ i ][ j ] = INF;
dst[ i ][ i ] = 0;
    }
  }
  void add_edge( int ui , int vi , int wi ){
    dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
    dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
  void shortest_path(){
    for( int k = 0 ; k < n ; k ++ )</pre>
       for( int i = 0 ; i < n ; i ++ )</pre>
         for( int j = 0 ; j < n ; j ++ )</pre>
           dst[ i ][ j ] = min( dst[ i ][ j ],
                  dst[ i ][ k ] + dst[ k ][ j ] );
  int solve( const vector<int>& ter ){
    int t = (int)ter.size();
for( int i = 0 ; i < ( 1 << t ) ; i ++ )</pre>
       for( int j = 0 ; j < n ; j ++ )</pre>
         dp[ i ][ j ] = INF;
    for( int i = 0 ; i < n ; i ++ )</pre>
       dp[0][i] = 0;
    for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
       if( msk == ( msk & (-msk) ) ){
         int who = __lg( msk );
for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];</pre>
         continue:
       for( int i = 0 ; i < n ; i ++ )</pre>
         for( int submsk = ( msk - 1 ) & msk ; submsk ;
                   submsk = (submsk - 1) \& msk)
              dp[ msk ][ i ] = min( dp[ msk ][ i ],
                                dp[ submsk ][ i ] +
                                 dp[ msk ^ submsk ][ i ] );
       for( int i = 0 ; i < n ; i ++ ){</pre>
         tdst[ i ] = INF;
```

### 3.12 Directed Minimum Spanning Tree

```
template <typename T> struct DMST {
   T g[maxn][maxn], fw[maxn];
   int n, fr[maxn];
   bool vis[maxn], inc[maxn];
   void clear() {
     for(int i = 0; i < maxn; ++i) {</pre>
       for(int j = 0; j < maxn; ++j) g[i][j] = inf;
vis[i] = inc[i] = false;</pre>
  }
   void addEdge(int u,int v,T w){g[u][v]=min(g[u][v],w)
   T operator()(int root, int _n) {
     n = _n; T ans = 0;
     if (dfs(root) != n) return -1;
     while (true) {
       for(int i = 1;i <= n;++i) fw[i] = inf, fr[i] = i;</pre>
       for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
         for (int j = 1; j <= n; ++j) {</pre>
           if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
             fw[i] = g[j][i]; fr[i] = j;
         }
       }
       int x = -1;
       for(int i = 1;i <= n;++i)if(i != root && !inc[i])</pre>
         int j = i, c = 0;
         while(j!=root && fr[j]!=i && c<=n) ++c, j=fr[j</pre>
         if (j == root || c > n) continue;
         else { x = i; break; }
       if (!~x) {
         for (int i = 1; i <= n; ++i)</pre>
           if (i != root && !inc[i]) ans += fw[i];
         return ans:
       int y = x;
       for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
       do {
         ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true
       } while (y != x);
       inc[x] = false;
       for (int k = 1; k <= n; ++k) if (vis[k]) {</pre>
         for (int j = 1; j <= n; ++j) if (!vis[j]) {</pre>
           if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
            if (g[j][k] < inf && g[j][k]-fw[k] < g[j][x])
              g[j][x] = g[j][k] - fw[k];
         }
       }
     }
     return ans;
   int dfs(int now) {
     int r = 1; vis[now] = true;
     for (int i = 1; i <= n; ++i)</pre>
       if (g[now][i] < inf && !vis[i]) r += dfs(i);</pre>
     return r;
  }
};
```

#### 3.13 Dominator Tree

```
\begin{tabular}{ll} \textbf{namespace} & \texttt{dominator} & \{ \end{tabular} \label{eq:continuous_parameter_space} \end{tabular}
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
  // vertices are numbered from 0 to n - 1
  fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
  fill(fa, fa + n, -1); fill(val, val + n, -1);
  fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
  fill(dom, dom + n, -1); tk = 0;
  for (int i = 0; i < n; ++i) {</pre>
    g[i].clear(); r[i].clear(); rdom[i].clear();
  }
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
  rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
  for (int u : g[x]) {
    if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
    r[dfn[u]].push_back(dfn[x]);
  }
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
  if (fa[x] == x) return c ? -1 : x;
  int p = find(fa[x], 1);
  if (p == -1) return c ? fa[x] : val[x];
  if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
  fa[x] = p;
  return c ? p : val[x];
vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
// p[i] = -2 if i is unreachable from s
  dfs(s);
  for (int i = tk - 1; i >= 0; --i) {
    for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)])
    if (i) rdom[sdom[i]].push_back(i);
    for (int &u : rdom[i]) {
      int p = find(u);
      if (sdom[p] == i) dom[u] = i;
      else dom[u] = p;
    if (i) merge(i, rp[i]);
  }
  vector < int > p(n, -2); p[s] = -1;
  for (int i = 1; i < tk; ++i)</pre>
    if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
  for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
  return p;
```

# 4 Matching & Flow

#### 4.1 Kuhn Munkres

```
class KM {
private:
  static constexpr lld INF = 1LL << 60;</pre>
  vector<lld> hl,hr,slk;
  vector<int> fl,fr,pre,qu;
  vector<vector<lld>> w;
  vector<bool> v1,vr;
  int n, q1, qr;
  bool check(int x) {
    if (vl[x] = true, fl[x] != -1)
      return vr[qu[qr++] = f1[x]] = true;
    while (x != -1) swap(x, fr[fl[x] = pre[x]]);
    return false;
  }
  void bfs(int s) {
    fill(slk.begin(), slk.end(), INF);
    fill(vl.begin(), vl.end(), false);
    fill(vr.begin(), vr.end(), false);
    ql = qr = 0;
    qu[qr++] = s;
    vr[s] = true;
    while (true) {
      11d d;
```

```
while (ql < qr) {
  for (int x = 0, y = qu[ql++]; x < n; ++x) {</pre>
           if(!vl[x] \&\& slk[x]>=(d=hl[x]+hr[y]-w[x][y]))
             if (pre[x] = y, d) slk[x] = d;
             else if (!check(x)) return;
        }
      }
      d = INF;
       for (int x = 0; x < n; ++x)
        if (!vl[x] && d > slk[x]) d = slk[x];
       for (int x = 0; x < n; ++x) {
        if (v1[x]) h1[x] += d;
         else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
       for (int x = 0; x < n; ++x)
         if (!vl[x] && !slk[x] && !check(x)) return;
    }
public:
  void init( int n_ ) {
    n = n_; qu.resize(n);
    fl.clear(); fl.resize(n, -1);
    fr.clear(); fr.resize(n, -1);
    hr.clear(); hr.resize(n); hl.resize(n);
    w.clear(); w.resize(n, vector<lld>(n));
    slk.resize(n); pre.resize(n);
    vl.resize(n); vr.resize(n);
  void set_edge( int u, int v, lld x ) { w[u][v] = x; }
  11d solve() {
    for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(w[i].begin(), w[i].end());
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    11d res = 0;
     for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
} km;
```

# 4.2 Bipartite Matching

```
class BipartiteMatching{
  private:
    vector<int> X[N], Y[N];
    int fX[N], fY[N], n;
    bitset<N> walked;
    bool dfs(int x){
      for(auto i:X[x]){
        if(walked[i])continue;
        walked[i]=1;
        if(fY[i]==-1||dfs(fY[i])){
           fY[i]=x;fX[x]=i;
          return 1:
        }
      }
      return 0;
  public:
    void init(int _n){
      for(int i=0;i<n;i++){</pre>
        X[i].clear();
        Y[i].clear();
        fX[i]=fY[i]=-1;
      walked.reset();
    void add_edge(int x, int y){
      X[x].push_back(y);
      Y[y].push_back(y);
    int solve(){
      int cnt = 0;
      for(int i=0;i<n;i++){</pre>
        walked.reset();
        if(dfs(i)) cnt++;
      // return how many pair matched
```

```
return cnt;
}
};
```

# 4.3 General Graph Matching

```
const int N = 514, E = (2e5) * 2;
struct Graph{
  int to[E],bro[E],head[N],e;
  int lnk[N],vis[N],stp,n;
  void init( int _n ){
    stp = 0; e = 1; n = _n;
    for( int i = 0 ; i <= n ; i ++ )</pre>
      head[i] = lnk[i] = vis[i] = 0;
  void add_edge(int u,int v){
    // 1-base
    to[e]=v,bro[e]=head[u],head[u]=e++;
    to[e]=u,bro[e]=head[v],head[v]=e++;
  bool dfs(int x){
    vis[x]=stp;
    for(int i=head[x];i;i=bro[i]){
      int v=to[i];
      if(!lnk[v]){
        lnk[x]=v, lnk[v]=x;
        return true;
      }else if(vis[lnk[v]]<stp){</pre>
        int w=lnk[v]:
        lnk[x]=v, lnk[v]=x, lnk[w]=0;
         if(dfs(w)) return true;
        lnk[w]=v, lnk[v]=w, lnk[x]=0;
      }
    return false;
  int solve(){
    int ans = 0;
    for(int i=1;i<=n;i++)</pre>
      if(not lnk[i]){
        stp++; ans += dfs(i);
    return ans;
} graph;
```

# 4.4 Minimum Weight Matching (Clique version)

```
struct Graph {
  // 0-base (Perfect Match)
  int n, edge[MXN][MXN];
  int match[MXN],dis[MXN],onstk[MXN];
  vector<int> stk;
  void init(int _n) {
    n = _n;
    for (int i=0; i<n; i++)</pre>
      for (int j=0; j<n; j++)</pre>
        edge[i][j] = 0;
  void set_edge(int u, int v, int w) {
    edge[u][v] = edge[v][u] = w;
  bool SPFA(int u){
    if (onstk[u]) return true;
    stk.PB(u);
    onstk[u] = 1;
    for (int v=0; v<n; v++){</pre>
      if (u != v && match[u] != v && !onstk[v]){
        int m = match[v];
        if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
          dis[m] = dis[u] - edge[v][m] + edge[u][v];
          onstk[v] = 1;
          stk.PB(v);
          if (SPFA(m)) return true;
          stk.pop_back();
          onstk[v] = 0;
     }
```

```
onstk[u] = 0;
     stk.pop_back();
     return false;
   int solve() {
     // find a match
     for (int i=0; i<n; i+=2){</pre>
      match[i] = i+1;
       match[i+1] = i;
     while (true){
       int found = 0;
       for (int i=0; i<n; i++)</pre>
         dis[i] = onstk[i] = 0;
       for (int i=0; i<n; i++){</pre>
         stk.clear();
         if (!onstk[i] && SPFA(i)){
           found = 1;
           while (SZ(stk)>=2){
             int u = stk.back(); stk.pop_back();
             int v = stk.back(); stk.pop_back();
             match[u] = v;
             match[v] = u;
         }
       if (!found) break;
     int ret = 0;
     for (int i=0; i<n; i++)</pre>
      ret += edge[i][match[i]];
     return ret>>1:
  }
} graph;
```

#### 4.5 Flow Models

- $\bullet$  Maximum/Minimum flow with lower/upper bound from s to t
  - 1. Construct super source  $\boldsymbol{S}$  and  $\sinh\,T$
  - 2. For each edge  $(x,y,l,u)\text{, connect }x\to y$  with capacity u-l
  - 3. For each vertex v, denote in(v) as the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds
  - 4. If in(v)>0 , connect  $S\to v$  with capacity in(v) , otherwise, connect  $v\to T$  with capacity -in(v)
    - To maximize, connect  $t \to s$  with capacity  $\infty$ , and let f be the maximum flow from S to T. If  $f \ne \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$  , where  $f_e$  corresponds to the flow on the graph
- Construct minimum vertex cover from maximum matching  ${\cal M}$  on bipartite graph (X,Y)
  - 1. Redirect every edge  $(y \to x \text{ if } (x,y) \in M$ ,  $x \to y \text{ otherwise})$
  - 2. DFS from unmatched vertices in X
  - 3.  $x \in X$  is chosen iff x is unvisited
  - 4.  $y \in Y$  is chosen iff y is visited
- Minimum cost cyclic flow
  - 1. Consruct super source  $\boldsymbol{S}$  and sink  $\boldsymbol{T}$
  - 2. For each edge (x,y,c), connect  $x\to y$  with (cost,cap)=(c,1) if c>0, otherwise connect  $y\to x$  with (cost,cap)=(-c,1)
  - 3. For each edge with c<0 , sum these cost as K , then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v)>0 , connect  $S\to v$  with (cost,cap)=(0,d(v))
  - 5. For each vertex v with d(v)<0 , connect  $v\to T$  with (cost,cap)=(0,-d(v))
  - 6. Flow from S to  $T\mbox{,}$  the answer is the cost of the flow C+K
- Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let  $\boldsymbol{K}$  be the sum of all weights

```
3. Connect source s \to v, v \in G with capacity K
4. For each edge (u,v,w) in G, connect u \to v and v \to u with capacity w
5. For v \in G, connect it with sink v \to t with capacity K + 2T - (\sum_{e \in E(v)} w(e)) - 2w(v)
6. T is a valid answer if the maximum flow f < K|V|
```

#### 4.6 Dinic

```
class Dinic{
private:
  using CapT = int64_t;
  struct Edge{
    int to, rev;
    CapT cap;
  int n, st, ed;
  vector<vector<Edge>> G;
  vector<int> lv;
  bool BFS(){
    fill(lv.begin(), lv.end(), -1);
    queue<int> bfs;
    bfs.push(st);
    lv[st] = 0;
    while(!bfs.empty()){
      int u = bfs.front(); bfs.pop();
      for(auto e: G[u]){
        if(e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
        lv[e.to] = lv[u] + 1;
        bfs.push(e.to);
    return (lv[ed]!=-1);
  CapT DFS(int u, CapT f){
    if(u == ed) return f;
    CapT ret = 0;
    for(auto& e: G[u]){
      if(e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
      CapT nf = DFS(e.to, min(f, e.cap));
      ret += nf; e.cap -= nf; f -= nf;
      G[e.to][e.rev].cap += nf;
      if(f == 0) return ret;
    if(ret == 0) lv[u] = -1;
    return ret:
public:
  void init(int n_, int st_, int ed_){
    n = n_, st = st_, ed = ed_;
    G.resize(n); lv.resize(n);
    fill(G.begin(), G.end(), vector<Edge>());
  void add_edge(int u, int v, CapT c){
    G[u].push_back({v, (int)G[v].size(), c});
    G[v].push_back({u, ((int)G[u].size())-1, 0});
  CapT max_flow(){
    CapT ret = 0;
    while(BFS()){
      CapT f = DFS(st, numeric_limits<CapT>::max());
      ret += f:
      if(f == 0) break;
    return ret;
} flow;
```

#### 4.7 Minimum Cost Maximum Flow

```
class MiniCostMaxiFlow{
  using CapT = int;
  using WeiT = int64_t;
  using PCW = pair<CapT,WeiT>;
  static constexpr CapT INF_CAP = 1 << 30;
  static constexpr WeiT INF_WEI = 1LL<<60;
private:
  struct Edge{
   int to, back;
   WeiT wei;</pre>
```

```
CapT cap:
     Edge() {}
     Edge(int a,int b,WeiT c,CapT d):
       to(a),back(b),wei(c),cap(d)
  };
  int ori, edd;
  vector<vector<Edge>> G;
  vector<int> fa, wh;
  vector<bool> inq;
  vector<WeiT> dis:
  PCW SPFA(){
     fill(inq.begin(),inq.end(),false);
     fill(dis.begin(),dis.end(),INF_WEI);
     queue<int> qq; qq.push(ori);
     dis[ori]=0;
     while(!qq.empty()){
       int u=qq.front();qq.pop();
       inq[u] = 0;
       for(int i=0;i<SZ(G[u]);++i){</pre>
         Edge e=G[u][i];
         int v=e.to;
         WeiT d=e.wei;
         if(e.cap<=0||dis[v]<=dis[u]+d)</pre>
           continue:
         dis[v]=dis[u]+d;
         fa[v]=u,wh[v]=i;
         if(inq[v]) continue;
         qq.push(v);
         inq[v]=1;
      }
     if(dis[edd]==INF WEI)
       return {-1,-1};
     CapT mw=INF_CAP;
     for(int i=edd;i!=ori;i=fa[i])
       mw=min(mw,G[fa[i]][wh[i]].cap);
     for (int i=edd;i!=ori;i=fa[i]){
       auto &eg=G[fa[i]][wh[i]];
       eg.cap-=mw;
       G[eg.to][eg.back].cap+=mw;
     return {mw,dis[edd]};
  }
public:
  void init(int a,int b,int n){
    ori=a,edd=b;
     G.clear();G.resize(n);
     fa.resize(n); wh.resize(n);
    inq.resize(n); dis.resize(n);
  void add_edge(int st,int ed,WeiT w,CapT c){
    G[st].emplace_back(ed,SZ(G[ed]),w,c);
     G[ed].emplace_back(st,SZ(G[st])-1,-w,0);
  PCW solve(){
    /* might modify to
     cc += ret.first * ret.second
     ww += ret.first * ret.second
     CapT cc=0; WeiT ww=0;
     while(true){
       PCW ret=SPFA();
       if(ret.first==-1) break;
       cc+=ret.first;
       ww+=ret.second;
     return {cc,ww};
} mcmf;
```

#### 4.8 Global Min-Cut

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];

void add_edge(int x, int y, int c) {
    w[x][y] += c;
    w[y][x] += c;
```

```
}
pair<int, int> phase(int n) {
    memset(v, false, sizeof(v));
    memset(g, 0, sizeof(g));
    int s = -1, t = -1;
    while (true) {
         int c = -1;
         for (int i = 0; i < n; ++i) {</pre>
              if (del[i] || v[i]) continue;
              if (c == -1 || g[i] > g[c]) c = i;
         if (c == -1) break;
         v[c] = true;
s = t, t = c;
         for (int i = 0; i < n; ++i) {</pre>
              if (del[i] || v[i]) continue;
              g[i] += w[c][i];
     return make_pair(s, t);
}
int mincut(int n) {
    int cut = 1e9;
    memset(del, false, sizeof(del));
for (int i = 0; i < n - 1; ++i) {</pre>
         int s, t; tie(s, t) = phase(n);
         del[t] = true;
         cut = min(cut, g[t]);
         for (int j = 0; j < n; ++j) {</pre>
             w[s][j] += w[t][j];
             w[j][s] += w[j][t];
         }
     return cut;
}
```

# 5 Math

#### 5.1 Prime Table

```
\begin{array}{c} 1002939109, 1020288887, 1028798297, 1038684299, \\ 1041211027, 1051762951, 1058585963, 1063020809, \\ 1147930723, 1172520109, 1183835981, 1187659051, \\ 1241251303, 1247184097, 1255940849, 1272759031, \\ 1287027493, 1288511629, 1294632499, 1312650799, \\ 1868732623, 1884198443, 1884616807, 1885059541, \\ 1909942399, 1914471137, 1923951707, 1925453197, \\ 1979612177, 1980446837, 1989761941, 2007826547, \\ 2008033571, 2011186739, 2039465081, 2039728567, \\ 2093735719, 2116097521, 2123852629, 2140170259, \\ 3148478261, 3153064147, 3176351071, 3187523093, \\ 3196772239, 3201312913, 3203063977, 3204840059, \\ 3210224309, 3213032591, 3217689851, 3218469083, \\ 3219857533, 3231880427, 3235951699, 3273767923, \\ 3276188869, 3277183181, 3282463507, 3285553889, \\ 3319309027, 3327005333, 3327574903, 3341387953, \\ 3373293941, 3380077549, 3380892997, 3381118801 \\ \end{array}
```

# 5.2 $\lfloor \frac{n}{i} \rfloor$ Enumeration

```
T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor
```

#### 5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 (mod n)
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
  if (y == 0) g=x,a=1,b=0;
  else
    exgcd(y,x%y,g,b,a),b-=(x/y)*a;
}
```

#### 5.4 Pollard Rho

```
// does not work when n is prime
// return any non-trivial factor
llu pollard_rho(llu n){
    static auto f=[](llu x,llu k,llu m){
        return add(k,mul(x,x,m),m);
    };
    if (!(n&1)) return 2;
    mt19937 rnd(120821011);
    while(true){
```

```
llu y=2,yy=y,x=rnd()%n,t=1;
for(llu sz=2;t==1;sz<<=1) {
   for(llu i=0;i<sz;++i){
      if(t!=1)break;
      yy=f(yy,x,n);
      t=gcd(yy>y?yy-y:y-yy,n);
   }
   y=yy;
}
if(t!=1&&t!=n) return t;
}
```

# 5.5 Pi Count (Linear Sieve)

```
static constexpr int N = 1000000 + 5;
11d pi[N];
vector<int> primes;
bool sieved[N];
lld cube_root(lld x){
  1ld s=cbrt(x-static_cast<long double>(0.1));
  while(s*s*s <= x) ++s;
  return s-1;
1ld square_root(lld x){
  1ld s=sqrt(x-static_cast<long double>(0.1));
  while(s*s <= x) ++s;
  return s-1;
void init(){
  primes.reserve(N);
  primes.push_back(1);
  for(int i=2;i<N;i++) {</pre>
    if(!sieved[i]) primes.push_back(i);
    pi[i] = !sieved[i] + pi[i-1];
    for(int p: primes) if(p > 1) {
   if(p * i >= N) break;
      sieved[p * i] = true;
      if(p % i == 0) break;
    }
  }
11d phi(lld m, lld n) {
  static constexpr int MM = 80000, NN = 500;
  static lld val[MM][NN];
  if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
  if(n == 0) return m;
  if(primes[n] >= m) return 1;
  lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
  if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
  return ret;
11d pi_count(11d);
11d P2(11d m, 11d n) {
  11d sm = square_root(m), ret = 0;
  for(lld i = n+1;primes[i]<=sm;i++)</pre>
    ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
  return ret;
1ld pi_count(lld m) {
  if(m < N) return pi[m];</pre>
  11d n = pi_count(cube_root(m));
  return phi(m, n) + n - 1 - P2(m, n);
```

#### 5.6 Range Sieve

```
const int MAX_SQRT_B = 50000;
const int MAX_L = 200000 + 5;

bool is_prime_small[MAX_SQRT_B];
bool is_prime[MAX_L];

void sieve(lld l, lld r){
    // [l, r)
    for(lld i=2;i*i<r;i++) is_prime_small[i] = true;
    for(lld i=1;i<r;i++) is_prime[i-1] = true;
    if(l==1) is_prime[0] = false;
    for(lld i=2;i*i<r;i++){
        if(!is_prime_small[i]) continue;
        for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;</pre>
```

#### 5.7 Miller Rabin

```
bool isprime(llu x){
  static llu magic[]={2,325,9375,28178,\
                     450775,9780504,1795265022};
  static auto witn=[](llu a,llu u,llu n,int t){
    a = mpow(a,u,n);
    if (!a)return 0;
    while(t--){
      1lu a2=mul(a,a,n);
      if(a2==1 && a!=1 && a!=n-1)
        return 1;
      a = a2;
    }
    return a!=1;
  if(x<2)return 0;</pre>
  if(!(x&1))return x==2;
  llu x1=x-1; int t=0;
  while(!(x1&1))x1>>=1,t++;
  for(llu m:magic)if(witn(m,x1,x,t))return 0;
  return 1;
}
```

#### 5.8 Inverse Element

```
// x's inverse mod k
long long GetInv(long long x, long long k){
    // k is prime: euler_(k)=k-1
    return qPow(x, euler_phi(k)-1);
}
// if you need [1, x] (most use: [1, k-1]
void solve(int x, long long k){
    inv[1] = 1;
    for(int i=2;i<x;i++)
        inv[i] = ((long long)(k - k/i) * inv[k % i]) % k;
}</pre>
```

## 5.9 Euler Phi Function

notprime[i\*j] = true;

```
extended euler:
   a^b mod p
   if gcd(a, p)==1: a^(b\%phi(p))
   elif b < phi(p): a^b mod p
   else a^(b%phi(p) + phi(p))
lld euler_phi(int x){
  11d r=1;
  for(int i=2;i*i<=x;++i){</pre>
    if(x%i==0){
      x/=i;
      r*=(i-1);
      while(x%i==0){
        x/=i;
        r*=i;
      }
    }
  if(x>1) r*=x-1;
  return r;
vector<int> primes;
bool notprime[N];
11d phi[N];
void euler_sieve(int n){
  for(int i=2;i<n;i++){</pre>
    if(!notprime[i]){
      primes.push_back(i);
      phi[i] = i-1;
    for(auto j: primes){
      if(i*j >= n) break;
```

```
phi[i*j] = phi[i] * phi[j];
if(i % j == 0){
    phi[i*j] = phi[i] * j;
    break;
}
}
}
}
```

#### 5.10 Gauss Elimination

#### 5.11 Fast Fourier Transform

```
polynomial multiply:
   DFT(a, len); DFT(b, len);
   for(int i=0;i<len;i++) c[i] = a[i]*b[i];</pre>
   iDFT(c, len);
   (len must be 2^k and = 2^k(max(a, b)))
   Hand written Cplx would be 2x faster
Cplx omega[2][N];
void init_omega(int n) {
  static constexpr llf PI=acos(-1);
  const llf arg=(PI+PI)/n;
  for(int i=0;i<n;++i)</pre>
    omega[0][i]={cos(arg*i),sin(arg*i)};
  for(int i=0;i<n;++i)</pre>
    omega[1][i]=conj(omega[0][i]);
void tran(Cplx arr[],int n,Cplx omg[]) {
  for(int i=0,j=0;i<n;++i){</pre>
    if(i>j)swap(arr[i],arr[j]);
    for(int l=n>>1;(j^=l)<l;l>>=1);
  for (int l=2;l<=n;l<<=1){</pre>
    int m=l>>1;
    for(auto p=arr;p!=arr+n;p+=1){
      for(int i=0;i<m;++i){</pre>
        Cplx t=omg[n/1*i]*p[m+i];
        p[m+i]=p[i]-t; p[i]+=t;
      }
    }
  }
void DFT(Cplx arr[],int n){tran(arr,n,omega[0]);}
void iDFT(Cplx arr[],int n){
  tran(arr,n,omega[1]);
  for(int i=0;i<n;++i) arr[i]/=n;</pre>
```

# 5.12 High Speed Linear Recurrence

```
#define mod 998244353
const int N=1000010;
int n,k,m,f[N],h[N],a[N],b[N],ib[N];
int pw(int x,int y){
  int re=1;
  if(y<0)y+=mod-1;
  while(y){
   if(y&1)re=(11)re*x%mod;
  y>>=1;x=(11)x*x%mod;
```

```
}
  return re;
void inc(int&x,int y){x+=y;if(x>=mod)x-=mod;}
namespace poly{
  const int G=3;
  int rev[N],L;
  void ntt(int*A,int len,int f){
    for(L=0;(1<<L)<len;++L);</pre>
    for(int i=0;i<len;++i){</pre>
      rev[i]=(rev[i>>1]>>1)|((i&1)<<(L-1));
      if(i<rev[i])swap(A[i],A[rev[i]]);</pre>
    for(int i=1;i<len;i<<=1){</pre>
      int wn=pw(G,f*(mod-1)/(i<<1));</pre>
      for(int j=0;j<len;j+=i<<1){</pre>
        int w=1;
        for(int k=0;k<i;++k,w=(11)w*wn%mod){</pre>
           int x=A[j+k],y=(11)w*A[j+k+i]%mod;
           A[j+k]=(x+y)\%mod,A[j+k+i]=(x-y+mod)\%mod;
      }
    if(!~f){
      int iv=pw(len,mod-2);
      for(int i=0;i<len;++i)A[i]=(11)A[i]*iv%mod;</pre>
  }
  void cls(int*A,int l,int r){
    for(int i=1;i<r;++i)A[i]=0;}</pre>
  void cpy(int*A,int*B,int 1){
    for(int i=0;i<1;++i)A[i]=B[i];}</pre>
  void inv(int*A,int*B,int 1){
    if(l==1){B[0]=pw(A[0],mod-2);return;}
    static int t[N];
    int len=l<<1;</pre>
    inv(A,B,l>>1);
    cpy(t,A,1);cls(t,1,len);
    ntt(t,len,1);ntt(B,len,1);
    for(int i=0;i<len;++i)</pre>
      B[i]=(11)B[i]*(2-(11)t[i]*B[i]%mod+mod)%mod;
    ntt(B,len,-1);cls(B,l,len);
  void pmod(int*A){
    static int t[N];
    int l=k+1,len=1;while(len<=(k<<1))len<<=1;</pre>
    cpy(t,A,(k<<1)+1);
    reverse(t,t+(k<<1)+1);
    cls(t,1,len);
    ntt(t,len,1);
    for(int i=0;i<len;++i)t[i]=(11)t[i]*ib[i]%mod;</pre>
    ntt(t,len,-1);
    cls(t,1,len);
    reverse(t,t+1);
    ntt(t,len,1);
    for(int i=0;i<len;++i)t[i]=(11)t[i]*b[i]%mod;</pre>
    ntt(t,len,-1);
    cls(t,1,len);
    for(int i=0;i<k;++i)A[i]=(A[i]-t[i]+mod)%mod;</pre>
    cls(A,k,len);
  void pow(int*A,int n){
    if(n==1){cls(A,0,k+1);A[1]=1;return;}
    pow(A,n>>1);
    int len=1; while(len<=(k<<1))len<<=1;</pre>
    ntt(A,len,1);
    for(int i=0;i<len;++i)A[i]=(11)A[i]*A[i]%mod;</pre>
    ntt(A,len,-1);
    pmod(A);
    if(n&1){
      for(int i=k;i;--i)A[i]=A[i-1];A[0]=0;
      pmod(A);
    }
 }
int main(){
 n=rd();k=rd();
  for(int i=1;i<=k;++i)f[i]=(mod+rd())%mod;</pre>
  for(int i=0;i<k;++i)h[i]=(mod+rd())%mod;</pre>
  for(int i=a[k]=b[k]=1;i<=k;++i)</pre>
    a[k-i]=b[k-i]=(mod-f[i])%mod;
  int len=1; while(len<=(k<<1))len<<=1;</pre>
```

```
reverse(a,a+k+1);
poly::inv(a,ib,len);
poly::cls(ib,k+1,len);
poly::ntt(b,len,1);
poly::ntt(ib,len,1);
poly::pow(a,n);
int ans=0;
for(int i=0;i<k;++i)inc(ans,(ll)a[i]*h[i]%mod);
printf("%d\n",ans);
return 0;
}</pre>
```

#### 5.13 Chinese Remainder

```
lld crt(lld ans[], lld pri[], int n){
  11d M = 1;
  for(int i=0;i<n;i++) M *= pri[i];</pre>
  11d ret = 0;
  for(int i=0;i<n;i++){</pre>
    lld inv = (gcd(M/pri[i], pri[i]).first + pri[i])%
    ret += (ans[i]*(M/pri[i])%M * inv)%M;
    ret %= M;
  return ret;
}
Another:
x = a1 \% m1
x = a2 \% m2
g = gcd(m1, m2)
assert((a1-a2)\%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
\theta \leftarrow x \leftarrow lcm(m1, m2)
```

### 5.14 Berlekamp Massey

```
// x: 1-base, p[]: 0-base
template < size_t N>
vector<llf> BM(llf x[N], size_t n){
  size_t f[N]={0},t=0;11f d[N];
  vector<llf> p[N];
  for(size_t i=1,b=0;i<=n;++i) {</pre>
     for(size_t j=0;j<p[t].size();++j)</pre>
    d[i]+=x[i-j-1]*p[t][j];
if(abs(d[i]-=x[i])<=EPS)continue;</pre>
    f[t]=i;if(!t){p[++t].resize(i);continue;}
    vector<llf> cur(i-f[b]-1);
    11f k=-d[i]/d[f[b]];cur.PB(-k);
     for(size_t j=0;j<p[b].size();j++)</pre>
      cur.PB(p[b][j]*k);
    if(cur.size()<p[t].size())cur.resize(p[t].size());</pre>
    for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];</pre>
    if(i-f[b]+p[b].size()>=p[t].size()) b=t;
    p[++t]=cur;
  return p[t];
```

#### 5.15 NTT

```
// Remember coefficient are mod P
/* p=a*2^n+1
   n
        2^n
                                а
                                      root
   16
        65536
                     65537
                                      3
                                1
                                      3 */
   20
        1048576
                     7340033
// (must be 2^k)
template < LL P, LL root, int MAXN >
struct NTT{
  static LL bigmod(LL a, LL b) {
    LL res = 1;
    for (LL bs = a; b; b >>= 1, bs = (bs * bs) % P)
      if(b&1) res=(res*bs)%P;
    return res;
  static LL inv(LL a, LL b) {
    if(a==1)return 1;
    return (((LL)(a-inv(b%a,a))*b+1)/a)%b;
```

```
LL omega[MAXN+1];
  NTT() {
    omega[0] = 1;
    LL r = bigmod(root, (P-1)/MAXN);
    for (int i=1; i<=MAXN; i++)</pre>
      omega[i] = (omega[i-1]*r)%P;
  // n must be 2^k
  void tran(int n, LL a[], bool inv_ntt=false){
    int basic = MAXN / n , theta = basic;
for (int m = n; m >= 2; m >>= 1) {
       int mh = m >> 1;
       for (int i = 0; i < mh; i++) {</pre>
         LL w = omega[i*theta%MAXN];
         for (int j = i; j < n; j += m) {</pre>
           int k = j + mh;
           LL x = a[j] - a[k];
           if (x < 0) x += P;
           a[j] += a[k];
           if (a[j] > P) a[j] -= P;
a[k] = (w * x) % P;
         }
      theta = (theta * 2) % MAXN;
    int i = 0;
    for (int j = 1; j < n - 1; j++) {</pre>
      for (int k = n >> 1; k > (i ^= k); k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
    if (inv_ntt) {
      LL ni = inv(n,P);
       reverse( a+1 , a+n );
      for (i = 0; i < n; i++)</pre>
         a[i] = (a[i] * ni) % P;
 }
};
const LL P=2013265921, root=31;
const int MAXN=4194304;
NTT<P, root, MAXN> ntt;
```

# 5.16 Polynomial Sqrt

```
const int mod = (119 << 23) + 1;</pre>
int inv_temp[400010];
void poly_inv(int *f, int *inv, int len) {
  int *inv_t = inv_temp, *g = inv;
  g[0] = get_inv(f[0]);
  for (int 1 = 2; 1 <= len; 1 <<= 1, swap(g, inv_t)) {</pre>
     for (int i = 0; i < 1; i++) {</pre>
       inv_t[i] = f[i];
       g[i + 1] = inv_t[i + 1] = 0;
     exec_ntt(inv_t, 1 << 1, 1);
     exec_ntt(g, 1 << 1, 1);
for (int i = 0; i < 2 * 1; i++)
      inv_t[i] = (ll)inv_t[i] * g[i] % mod;
     for (int i = 0; i < 2 * 1; i++) {
       if (inv_t[i])
         inv_t[i] = mod - inv_t[i];
       inv_t[i] += 2, inv_t[i] %= mod;
     for (int i = 0; i < 2 * 1; i++)
       inv_t[i] = (ll)inv_t[i] * g[i] % mod;
     exec_ntt(inv_t, l << 1, -1);
for (int i = 0; i < l; i++)
       inv_t[i + 1] = 0;
  for (int i = 0; i < len; i++)</pre>
    inv[i] = g[i];
int sqrt_temp[400010], inv_t[400010];
void poly_sqrt(int *f, int *sqrt_pol, int len) {
   int *g = sqrt_pol, *t = sqrt_temp, inv2 = get_inv(2);
}
  g[0] = 1;
  for (int 1 = 2; 1 <= len; 1 <<= 1, swap(g, t)) {</pre>
     for (int i = 0; i < 1; i++)</pre>
       t[i] = f[i], t[i + 1] = g[i + 1] = inv_t[i] = 0;
     poly_inv(g, inv_t, 1);
```

```
for (int i = 1; i < 2 * 1; i++)
      inv_t[i] = 0;
     exec_ntt(g, l << 1, 1);
    exec_ntt(inv_t, 1 << 1, 1);
     exec_ntt(t, 1 << 1, 1);
    for (int i = 0; i < (1 << 1); i++)</pre>
      t[i]=(ll)inv2*(g[i]+(ll)t[i]*inv_t[i] % mod)%mod;
     exec_ntt(t, 1 << 1, -1);
    for (int i = 0; i < 1; i++)</pre>
       t[i + 1] = 0;
  for (int i = 0; i < len; i++)</pre>
    sqrt_pol[i] = g[i];
int c[400010], inv[400010], sqrt_pol[400010];
int main(){
  int n, m, x;
scanf("%d%d", &n, &m);
  for (int i = 0; i < n; i++)</pre>
    scanf("%d", &x);
    if (x <= m)
      c[x] = mod - 4;
  c[0]++, c[0] \% = mod;
  int len = 1;
  while (len <= m)len <<= 1;</pre>
  poly_sqrt(c, sqrt_pol, len);
  sqrt_pol[0]++, sqrt_pol[0] %= mod;
  poly_inv(sqrt_pol, inv, len);
for (int i = 1; i <= m; i++)</pre>
    printf("%d\n", (inv[i] + inv[i]) % mod);
  puts("");
  return 0;
```

# 5.17 Polynomial Division

```
VI inverse(const VI &v, int n) {
  VI q(1, fpow(v[0], mod - 2));
  for (int i = 2; i <= n; i <<= 1) {</pre>
    VI fv(v.begin(), v.begin() + i);
    VI fq(q.begin(), q.end());
    fv.resize(2 * i), fq.resize(2 * i);
    ntt(fq, 2 * i), ntt(fv, 2 * i);
for (int j = 0; j < 2 * i; ++j)
      fv[j] = fv[j]*1ll*fq[j]%mod*fq[j]%mod;
    intt(fv, 2 * i);
    VI res(i);
    for (int j = 0; j < i; ++j) {</pre>
      res[j] = mod - fv[j];
      if (j < (i>>1)) (res[j] += 2*q[j]%mod) %= mod;
    q = res;
  }
  return q;
VI divide(const VI &a, const VI &b) {
  // leading zero should be trimmed
  int n = (int)a.size(), m = (int)b.size();
  int k = 2;
  while (k < n - m + 1) k <<= 1;
  VI ra(k), rb(k);
  for (int i = 0; i < min(n, k); ++i) ra[i] = a[n-i-1];</pre>
  for (int i = 0; i < min(m, k); ++i) rb[i] = b[m-i-1];</pre>
  VI rbi = inverse(rb, k):
  VI res = convolution(rbi, ra);
  res.resize(n - m + 1);
  reverse(res.begin(), res.end());
  return res;
```

### 5.18 FWT

```
/* xor convolution:

* x = (x0,x1) , y = (y0,y1)

* z = ( x0y0 + x1y1 , x0y1 + x1y0 )

* =>

* x' = ( x0+x1 , x0-x1 ) , y' = ( y0+y1 , y0-y1 )

* z' = ( ( x0+x1 )( y0+y1 ) , ( x0-x1 )( y0-y1 ) )

* z = (1/2) * z''
```

```
* or convolution:
  x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
 * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
const LL MOD = 1e9+7;
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
  for( int d = 1 ; d < N ; d <<= 1 ) {</pre>
    int d2 = d << 1;
    for( int s = 0 ; s < N ; s += d2 )</pre>
       for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
  LL ta = x[ i ] , tb = x[ j ];</pre>
         x[i] = ta+tb;
         x[ j ] = ta-tb;
         if( x[ i ] >= MOD ) x[ i ] -= MOD;
         if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
  if( inv )
    for( int i = 0 ; i < N ; i++ ) {</pre>
      x[ i ] *= inv( N, MOD );
      x[ i ] %= MOD;
}
```

### 5.19 DiscreteLog

```
// Baby-step Giant-step Algorithm
11d BSGS(11d P, 11d B, 11d N) {
  // find B^L = N mod P
  unordered_map<lld, int> R;
  11d sq = (11d)sqrt(P);
  11d t = 1;
  for (int i = 0; i < sq; i++) {</pre>
   if (t == N) return i;
    if (!R.count(t)) R[t] = i;
   t = (t * B) \% P;
 11d f = inverse(t, P);
 for(int i=0;i<=sq+1;i++) {</pre>
   if (R.count(N))
     return i * sq + R[N];
   N = (N * f) % P;
  return -1;
```

#### 5.20 Quadratic residue

```
struct Status{
 11 x,y;
11 w;
Status mult(const Status& a,const Status& b,ll mod){
 res.x=(a.x*b.x+a.y*b.y%mod*w)%mod;
 res.y=(a.x*b.y+a.y*b.x)%mod;
 return res;
inline Status qpow(Status _base,11 _pow,11 _mod){
 Status res = \{1, 0\};
 while(_pow>0){
   if(_pow&1) res=mult(res,_base,_mod);
    _base=mult(_base,_base,_mod);
   _pow>>=1;
 }
 return res;
inline 11 check(11 x,11 p){
 return qpow_mod(x,(p-1)>>1,p);
inline ll get_root(ll n,ll p){
 if(p==2) return 1;
 if(check(n,p)==p-1) return -1;
 ll a;
 while(true){
    a=rand()%p;
    w=((a*a-n)%p+p)%p;
   if(check(w,p)==p-1) break;
 Status res = {a, 1}
 res=qpow(res,(p+1)>>1,p);
```

```
5.21 De-Bruijn
```

return res.x:

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
  if (t > n) {
    if (n % p == 0)
      for (int i = 1; i <= p; ++i)</pre>
        res[sz++] = aux[i];
  } else +
    aux[t] = aux[t - p];
    db(t + 1, p, n, k);
for (int i = aux[t - p] + 1; i < k; ++i) {
      aux[t] = i;
      db(t + 1, t, n, k);
  }
int de_bruijn(int k, int n) {
  // return cyclic string of len k^n s.t. every string
  // of len n using k char appears as a substring.
  if (k == 1) {
    res[0] = 0;
    return 1;
  for (int i = 0; i < k * n; i++) aux[i] = 0;</pre>
  sz = 0;
  db(1, 1, n, k);
  return sz;
```

# 5.22 Simplex Construction

Standard form: maximize  $\sum_{1\leq i\leq n}c_ix_i$  such that for all  $1\leq j\leq m$ ,  $\sum_{1\leq i\leq n}A_{ji}x_i\leq b_j.$ and  $x_i\geq 0$  for all  $1\leq i\leq n.$ 

- 1. In case of minimization, let  $c_i^\prime = -c_i$
- 2.  $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \rightarrow \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- 3.  $\sum_{1 \leq i \leq n} A_{ji} x_i = b_j$ 
  - $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$
  - $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

#### 5.23 Simplex

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return vector<double>(n, -inf) if the solution doesn
// return vector<double>(n, +inf) if the solution is
    unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
  double inv = 1.0 / d[r][s];
  for (int i = 0; i < m + 2; ++i) {</pre>
    for (int j = 0; j < n + 2; ++j) {
  if (i != r && j != s)</pre>
        d[i][j] -= d[r][j] * d[i][s] * inv;
    }
  for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;</pre>
  for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
  d[r][s] = inv;
  swap(p[r], q[s]);
bool phase(int z) {
  int x = m + z;
  while (true) {
    int s = -1;
```

```
for (int i = 0; i <= n; ++i) {</pre>
      if (!z && q[i] == -1) continue;
       if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
    if (d[x][s] > -eps) return true;
    int r = -1;
    for (int i = 0; i < m; ++i) {</pre>
      if (d[i][s] < eps) continue;</pre>
      if (r == -1 || \
         d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
    if (r == -1) return false;
    pivot(r, s);
  }
VD solve(const VVD &a, const VD &b, const VD &c) {
  m = b.size(), n = c.size();
  d = VVD(m + 2, VD(n + 2));
  for (int i = 0; i < m; ++i) {</pre>
    for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
  p.resize(m), q.resize(n + 1);
  for (int i = 0; i < m; ++i)</pre>
    p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
  for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];</pre>
  q[n] = -1, d[m + 1][n] = 1;
  int r = 0;
  for (int i = 1; i < m; ++i)</pre>
    if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
  if (d[r][n + 1] < -eps) {</pre>
    pivot(r, n);
    if (!phase(1) || d[m + 1][n + 1] < -eps)
    return VD(n, -inf);
for (int i = 0; i < m; ++i) if (p[i] == -1) {</pre>
      int s = min_element(d[i].begin(), d[i].end() - 1)
                - d[i].begin();
      pivot(i, s);
    }
  if (!phase(0)) return VD(n, inf);
  VD x(n);
  for (int i = 0; i < m; ++i)</pre>
    if (p[i] < n) x[p[i]] = d[i][n + 1];</pre>
  return x;
}}
```

# 6 Geometry

#### 6.1 Point Class

```
template<typename T>
struct Point{
 typedef long double llf;
  static constexpr llf EPS = 1e-8;
  T x, y;
 Point(T _=0, T __=0): x(_), y(__){}
  template<typename T2>
    Point(const Point<T2>& a): x(a.x), y(a.y){}
  inline llf theta() const {
    return atan2((11f)y, (11f)x);
 inline llf dis() const {
    return hypot((llf)x, (llf)y);
 inline llf dis(const Point& o) const {
    return hypot((llf)(x-o.x), (llf)(y-o.y));
  Point operator-(const Point& o) const {
    return Point(x-o.x, y-o.y);
  Point operator+(const Point& o) const {
   return Point(x+o.x, y+o.y);
  Point operator*(const T& k) const {
    return Point(x*k, y*k);
  Point operator/(const T& k) const {
   return Point(x/k, y/k);
 Point operator-() const {
    return Point(-x, -y);
```

```
Point rot90() const {
    return Point(-y, x);
  template<typename T2>
  bool in(const Circle<T2>& a) const {
    /* Add struct Circle at top */
    return a.o.dis(*this)+EPS <= a.r;</pre>
  bool equal(const Point& o, true_type) const {
    return fabs(x-o.x) \langle EPS and fabs(y-o.y) \langle EPS;
  bool equal(const Point& o, false_type) const {
    return tie(x, y) == tie(o.x, o.y);
  bool operator == (const Point& o) const {
    return equal(o, is_floating_point<T>());
  bool operator!=(const Point& o) const {
    return !(*this == 0);
  bool operator<(const Point& o) const {</pre>
    return theta() < o.theta();</pre>
    // sort like what pairs did
    // if (is_floating_point<T>())
         return fabs(x-o.x)<EPS?y<o.y:x<o.x;
    // else return tie(x, y) < tie(o.x, o.y);
  friend inline T cross(const Point& a, const Point& b)
    return a.x*b.y - b.x*a.y;
  friend inline T dot(const Point& a, const Point &b){
    return a.x*b.x + a.y*b.y;
  friend ostream& operator<<(ostream& ss, const Point&</pre>
    ss<<"("<<o.x<<", "<<o.y<<")";
    return ss;
};
```

### 6.2 Circle Class

#### 6.3 Segment Class

```
const Point<long double> INF_P(-1e20, 1e20);
const Point<long double> NOT_EXIST(1e20, 1e-20);
template<typename T>
struct Line{
  static constexpr long double EPS = 1e-8;
  // ax+by+c = 0
  T a, b, c;
  Line(): a(0), b(1), c(0){}
                , T
          , T
                        _): a(_), b(__), c(___){
  Line(T
    assert(fabs(a)>EPS or fabs(b)>EPS);
  template<typename T2>
    Line(const Line\langle T2 \rangle \& x): a(x.a), b(x.b), c(x.c){}
  typedef Point<long double> Pt;
  bool equal(const Line& o, true_type) const {
    return fabs(a-o.a) < EPS and fabs(b-o.b) < EPS and
        fabs(c-o.b) < EPS;</pre>
```

```
bool equal(const Line& o, false_type) const {
    return a==o.a and b==o.b and c==o.c;
  bool operator==(const Line& o) const {
    return equal(o, is_floating_point<T>());
  bool operator!=(const Line& o) const {
    return !(*this == 0);
  friend inline bool on_line__(const Point<T>& p, const
       Line& 1, true_type){
    return fabs(1.a*p.x + 1.b*p.y + 1.c) < EPS;</pre>
  friend inline bool on_line__(const Point<T>& p, const
       Line& 1, false_type){
    return 1.a*p.x + 1.b*p.y + 1.c == 0;
  friend inline bool on_line(const Point<T>&p, const
      Line& 1){
    return on_line__(p, 1, is_floating_point<T>());
  friend inline bool is_parallel__(const Line& x, const
       Line& y, true_type){
    return fabs(x.a*y.b - x.b*y.a) < EPS;</pre>
  friend inline bool is_parallel__(const Line& x, const
       Line& y, false_type){
    return x.a*y.b == x.b*y.a;
  friend inline bool is_parallel(const Line& x, const
      Line& y){
    return is_parallel__(x, y, is_floating_point<T>());
  friend inline Pt get_inter(const Line& x, const Line&
      y){
    typedef long double llf;
    if(x==y) return INF_P;
    if(is_parallel(x, y)) return NOT_EXIST;
    llf delta = x.a*y.b - x.b*y.a;
    llf delta_x = x.b*y.c - x.c*y.b;
    11f delta y = x.c*y.a - x.a*y.c;
    return Pt(delta_x / delta, delta_y / delta);
  friend ostream& operator<<(ostream& ss, const Line& o</pre>
    ss<<o.a<<"x+"<<o.b<<"y+"<<o.c<<"=0";
    return ss;
 }
template<typename T>
inline Line<T> get_line(const Point<T>& a, const Point<</pre>
    T>& b){}
  return Line<T>(a.y-b.y, b.x-a.x, (b.y-a.y)*a.x-(b.x-a
      .x)*a.y);
```

# 6.4 Line Class

```
const Point<long double> INF_P(-1e20, 1e20);
const Point<long double> NOT_EXIST(1e20, 1e-20);
template<typename T>
struct Line{
 static constexpr long double EPS = 1e-8;
  // ax+by+c = 0
 T a, b, c;
 Line(): a(0), b(1), c(0){}
 Line(T _, T
                 Т
                       _): a(_), b(_
                                     _), c(___){
    assert(fabs(a)>EPS or fabs(b)>EPS);
 template<typename T2>
   Line(const Line\langle T2 \rangle \& x): a(x.a), b(x.b), c(x.c){}
  typedef Point<long double> Pt;
 bool equal(const Line& o, true_type) const {
    return fabs(a-o.a) < EPS and fabs(b-o.b) < EPS and
        fabs(c-o.b) < EPS;
  bool equal(const Line& o, false_type) const {
   return a==o.a and b==o.b and c==o.c;
  bool operator==(const Line& o) const {
    return equal(o, is_floating_point<T>());
```

```
bool operator!=(const Line& o) const {
    return !(*this == 0);
  friend inline bool on_line__(const Point<T>& p, const
       Line& 1, true_type){
    return fabs(1.a*p.x + 1.b*p.y + 1.c) < EPS;</pre>
  friend inline bool on_line__(const Point<T>& p, const
       Line& 1, false_type){
    return 1.a*p.x + 1.b*p.y + 1.c == 0;
  friend inline bool on_line(const Point<T>&p, const
      Line& 1){
    return on_line__(p, l, is_floating_point<T>());
  friend inline bool is_parallel__(const Line& x, const
       Line& y, true_type){
    return fabs(x.a*y.b - x.b*y.a) < EPS;</pre>
  friend inline bool is_parallel__(const Line& x, const
       Line& y, false_type){
    return x.a*y.b == x.b*y.a;
  friend inline bool is_parallel(const Line& x, const
      Line& y){
    return is_parallel__(x, y, is_floating_point<T>());
  friend inline Pt get_inter(const Line& x, const Line&
    typedef long double llf;
    if(x==y) return INF_P;
    if(is_parallel(x, y)) return NOT_EXIST;
    llf delta = x.a*y.b - x.b*y.a;
    llf delta_x = x.b*y.c - x.c*y.b;
    11f delta_y = x.c*y.a - x.a*y.c;
    return Pt(delta_x / delta, delta_y / delta);
  friend ostream& operator<<(ostream& ss, const Line& o</pre>
    ss<<o.a<<"x+"<<o.b<<"y+"<<o.c<<"=0";
    return ss;
};
template<typename T>
inline Line<T> get_line(const Point<T>& a, const Point<</pre>
    T>& b){
  return Line<T>(a.y-b.y, b.x-a.x, (b.y-a.y)*a.x-(b.x-a
      .x)*a.y);
```

# 6.5 Triangle Circumcentre

```
template < typename T >
Circle < llf > get_circum(const Point < T > & a, const Point < T > & b, const Point < T > & c) {
    llf a1 = a.x-b.x;
    llf b1 = a.y-b.y;
    llf c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
    llf a2 = a.x-c.x;
    llf b2 = a.y-c.y;
    llf c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;

    Circle < llf > cc;
    cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
    cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
    cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
    return cc;
}
```

## 6.6 2D Convex Hull

```
template<typename T>
class ConvexHull_2D{
private:
   typedef Point<T> PT;
   vector<PT> dots;
   struct myhash{
    uint64_t operator()(const PT& a) const {
      uint64_t xx=0, yy=0;
      memcpy(&xx, &a.x, sizeof(a.x));
   }
}
```

```
memcpy(&yy, &a.y, sizeof(a.y));
      uint64_t ret = xx*17+yy*31;
      ret = (ret ^ (ret >> 16))*0x9E3779B1;
      ret = (ret ^ (ret >> 13))*0xC2B2AE35;
      ret = ret ^ xx;
      return (ret ^ (ret << 3)) * yy;</pre>
  };
  unordered_set<PT, myhash> in_hull;
public:
  inline void init(){in_hull.clear();dots.clear();}
  void insert(const PT& x){dots.PB(x);}
  void solve(){
    sort(ALL(dots), [](const PT& a, const PT& b){
      return tie(a.x, a.y) < tie(b.x, b.y);</pre>
    vector<PT> stk(SZ(dots)<<1);</pre>
    int top = 0;
    for(auto p: dots){
      while(top >= 2 && cross(p-stk[top-2], stk[top-1]-
          stk[top-2]) <= 0)
        top --;
      stk[top++] = p;
    for(int i=SZ(dots)-2, t = top+1;i>=0;i--){
      while(top >= t && cross(dots[i]-stk[top-2], stk[
          top-1]-stk[top-2]) <= 0)
        top --:
      stk[top++] = dots[i];
    stk.resize(top-1);
    swap(stk, dots);
    for(auto i: stk) in_hull.insert(i);
  vector<PT> get(){return dots;}
  inline bool in_it(const PT& x){
    return in_hull.find(x)!=in_hull.end();
  }
};
```

#### 6.7 2D Farthest Pair

#### 6.8 2D Closest Pair

```
struct Pt{
  11f x, y, d;
} arr[N];
inline llf dis(Pt a, Pt b){
  return hypot(a.x-b.x, a.y-b.y);
11f solve(){
  int cur = rand() % n;
  for(int i=0;i<n;i++) arr[i].d = dis(arr[cur], arr[i])</pre>
  sort(arr, arr+n, [](Pt a, Pt b){return a.d < b.d;});</pre>
  11f ans = 1e50;
  for(int i=0;i<n;i++){</pre>
    for(int j=i+1;j<n;j++){</pre>
      if(arr[j].d - arr[i].d > ans) break;
      ans = min(ans, dis(arr[i], arr[j]));
    }
  return ans;
```

# 6.9 SimulateAnnealing

```
llf anneal() {
  mt19937 rnd_engine( seed );
  uniform_real_distribution< llf > rnd( 0, 1 );
  const llf dT = 0.001;

// Argument p

llf S_cur = calc( p ), S_best = S_cur;
  for ( llf T = 2000 ; T > EPS ; T -= dT ) {
      // Modify p to p_prime
      const llf S_prime = calc( p_prime );
      const llf delta_c = S_prime - S_cur;
      llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
      if ( rnd( rnd_engine ) <= prob )
            S_cur = S_prime, p = p_prime;
      if ( S_prime < S_best )
            S_best = S_prime, p_best = p_prime;
    }
    return S_best;
}</pre>
```

#### 6.10 Half Plane Intersection

```
inline int dcmp ( double x ) {
  if( fabs( x ) < eps ) return 0;</pre>
  return x > 0 ? 1 : -1;
struct Line {
  Point st, ed;
  double ang;
  Line( Point _st=Point(), Point _ed=Point()):
   st(_st),ed(_ed),ang(atan2(_ed.y-_st.y,_ed.x-_st.x))
       {}
  inline bool operator< ( const Line& rhs ) const {</pre>
    if(dcmp( ang - rhs.ang ) != 0) return ang < rhs.ang</pre>
    return dcmp( cross( st, ed, rhs.st ) ) < 0;</pre>
};
// cross(pt, line.ed-line.st) >= 0 <-> pt in half plane
vector< Line > lines;
deque< Line > que;
deque< Point > pt;
double HPI() {
  sort( lines.begin(), lines.end() );
  que.clear();
  pt.clear();
  que.push_back( lines[ 0 ] );
  for ( int i = 1 ; i < (int)lines.size() ; i ++ ) {</pre>
    if(!dcmp(lines[i].ang - lines[i - 1].ang)) continue
    while ( pt.size() > 0 &&
     dcmp(cross(lines[i].st,lines[i].ed,pt.back()))<0){</pre>
      pt.pop_back();
      que.pop_back();
    while ( pt.size() > 0 &&
     dcmp(cross(lines[i].st,lines[i].ed,pt.front()))<0)</pre>
      pt.pop_front();
      que.pop_front();
    pt.push_back( get_point( que.back(), lines[ i ] ) )
    que.push_back( lines[ i ] );
  while ( pt.size() > 0 &&
   dcmp(cross(que[0].st, que[0].ed, pt.back()))<0){</pre>
    que.pop_back();
    pt.pop_back();
  while ( pt.size() > 0 &&
   dcmp(cross(que.back().st, que.back().ed, pt[0]))<0)</pre>
    que.pop_front();
    pt.pop_front();
  pt.push_back( get_point( que.front(), que.back() ) );
  vector< Point > conv;
  for ( int i = 0 ; i < (int)pt.size() ; i ++ )</pre>
    conv.push_back( pt[ i ] );
  double ret = 0;
  for ( int i = 1 ; i + 1 < (int)conv.size() ; i ++ )</pre>
```

```
ret += abs(cross(conv[0], conv[i], conv[i + 1]));
return ret / 2.0;
}
```

## 6.11 Ternary Search on Integer

```
int TernarySearch(int 1, int r) {
    // max value @ (l, r]
    while (r - 1 > 1){
        int m = (1 + r)>>1;
        if (f(m) > f(m + 1)) r = m;
        else l = m;
    }
    return l+1;
}
```

### 6.12 Minimum Covering Circle

```
template<typename T>
Circle<llf> MinCircleCover(const vector<Point<T>>& pts)
  random_shuffle(ALL(pts));
  Circle<llf> c = \{pts[0], 0\};
  int n = SZ(pts);
  for(int i=0;i<n;i++){</pre>
    if(pts[i].in(c)) continue;
    c = {pts[i], 0};
    for(int j=0;j<i;j++){</pre>
      if(pts[j].in(c)) continue;
      c.o = (pts[i] + pts[j]) / 2;
      c.r = pts[i].dis(c.o);
      for(int k=0;k<j;k++){</pre>
        if(pts[k].in(c)) continue;
        c = get_circum(pts[i], pts[j], pts[k]);
    }
  }
  return c;
}
```

#### 6.13 KDTree (Nearest Point)

```
const int MXN = 100005;
struct KDTree {
  struct Node {
    int x,y,x1,y1,x2,y2;
    int id,f;
Node *L, *R;
  } tree[MXN], *root;
  int n;
  LL dis2(int x1, int y1, int x2, int y2) {
    LL dx = x1-x2, dy = y1-y2;
    return dx*dx+dy*dy;
  static bool cmpx(Node& a, Node& b){ return a.x<b.x; }</pre>
  static bool cmpy(Node& a, Node& b){ return a.y<b.y; }</pre>
  void init(vector<pair<int,int>> ip) {
    n = ip.size();
    for (int i=0; i<n; i++) {</pre>
      tree[i].id = i;
      tree[i].x = ip[i].first;
      tree[i].y = ip[i].second;
    root = build_tree(0, n-1, 0);
  Node* build_tree(int L, int R, int dep) {
    if (L>R) return nullptr;
    int M = (L+R)/2; tree[M].f = dep%2;
    nth_element(tree+L,tree+M,tree+R+1,dep%2?cmpy:cmpx)
    tree[M].x1 = tree[M].x2 = tree[M].x;
    tree[M].y1 = tree[M].y2 = tree[M].y;
    tree[M].L = build_tree(L, M-1, dep+1);
    if (tree[M].L) {
      tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
      tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
    tree[M].R = build_tree(M+1, R, dep+1);
```

```
if (tree[M].R) {
      tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
      tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
      tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
    return tree+M;
  int touch(Node* r, int x, int y, LL d2){
    LL dis = sqrt(d2)+1;
    if (x<r->x1-dis || x>r->x2+dis ||
        y<r->y1-dis || y>r->y2+dis)
      return 0;
    return 1;
  }
  void nearest(Node* r, int x, int y,int &mID, LL &md2)
    if (!r || !touch(r, x, y, md2)) return;
    LL d2 = dis2(r->x, r->y, x, y);
    if (d2 < md2 \mid | (d2 == md2 && mID < r->id)) {
      mID = r -> id;
      md2 = d2;
    // search order depends on split dim
    if ((r->f == 0 && x < r->x) | |
        (r->f == 1 \&\& y < r->y))
      nearest(r->L, x, y, mID, md2);
      nearest(r->R, x, y, mID, md2);
    } else {
      nearest(r->R, x, y, mID, md2);
      nearest(r->L, x, y, mID, md2);
  int query(int x, int y) {
    int id = 1029384756;
    LL d2 = 102938475612345678LL;
    nearest(root, x, y, id, d2);
    return id;
  }
}tree;
```

# 7 Stringology

## 7.1 Hash

```
class Hash{
private:
  const int p = 127, q = 1208220623;
  int sz, prefix[N], power[N];
  inline int add(int x, int y){return x+y>=q?x+y-q:x+y
  inline int sub(int x, int y){return x-y<0?x-y+q:x-y;}</pre>
  inline int mul(int x, int y){return 1LL*x*y%q;}
public:
  void init(const string &x){
     sz = x.size();prefix[0]=0;power[0]=1;
     for(int i=1;i<=sz;i++)</pre>
       prefix[i]=add(mul(prefix[i-1], p), x[i-1]);
    for(int i=1;i<=sz;i++) power[i]=mul(power[i-1], p);</pre>
  int query(int 1, int r){
    // 1-base (l, r]
     return sub(prefix[r], mul(prefix[l], power[r-l]));
};
```

# 7.2 Suffix Array

```
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
int x[maxn], p[maxn], q[maxn * 2];
// sa[i]: sa[i]-th suffix is the \
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix \
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *sa, int *c, int n, int z) {
  memset(sa, 0, sizeof(int) * n);
  memcpy(x, c, sizeof(int) * z);
```

```
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
  memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
    if (sa[i] && !t[sa[i] - 1])
       sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  memcpy(x, c, sizeof(int) * z);
  for (int i = n - 1; i >= 0; --i)
    if (sa[i] && t[sa[i] - 1])
       sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q,
 bool *t, int *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
  memset(c, 0, sizeof(int) * z);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];</pre>
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
     return;
  for (int i = n - 2; i >= 0; --i)
    t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  pre(sa, c, n, z);
  for (int i = 1; i <= n - 1; ++i)</pre>
    if (t[i] && !t[i - 1])
       sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) {
    if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
    bool neq = last < 0 || \</pre>
      memcmp(s + sa[i], s + last,
  (p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
    ns[q[last = sa[i]]] = nmxz += neq;
  }}
  sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
  pre(sa, c, n, z);
  for (int i = nn - 1; i >= 0; --i)
    sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
void build(const string &s) {
  for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
  _s[(int)s.size()] = 0; // s shouldn't contain 0
  sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];</pre>
  for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
  int ind = 0; hi[0] = 0;
  for (int i = 0; i < (int)s.size(); ++i) {</pre>
    if (!rev[i]) {
       ind = 0;
       continue;
    while (i + ind < (int)s.size() && \</pre>
      s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
    hi[rev[i]] = ind ? ind-- : 0;
}}
```

# 7.3 Aho-Corasick Algorithm

```
class AhoCorasick{
 private:
    static constexpr int Z = 26;
    struct node{
     node *nxt[ Z ], *fail;
      vector< int > data;
     node(): fail( nullptr ) {
        memset( nxt, 0, sizeof( nxt ) );
        data.clear();
    } *rt;
    inline int Idx( char c ) { return c - 'a'; }
    void init() { rt = new node(); }
    void add( const string& s, int d ) {
     node* cur = rt;
      for ( auto c : s ) {
        if ( not cur->nxt[ Idx( c ) ] )
          cur->nxt[ Idx( c ) ] = new node();
```

```
cur = cur->nxt[ Idx( c ) ];
       cur->data.push_back( d );
     }
     void compile() {
       vector< node* > bfs;
       size_t ptr = 0;
       for ( int i = 0 ; i < Z ; ++ i ) {</pre>
         if ( not rt->nxt[ i ] )
           continue;
          rt->nxt[ i ]->fail = rt;
         bfs.push_back( rt->nxt[ i ] );
       while ( ptr < bfs.size() ) {</pre>
         node* u = bfs[ ptr ++ ];
for ( int i = 0 ; i < Z ; ++ i ) {
            if ( not u->nxt[ i ] )
             continue;
            node* u_f = u->fail;
            while ( u_f ) {
              if ( not u_f->nxt[ i ] ) {
                u_f = u_f->fail; continue;
              u->nxt[ i ]->fail = u_f->nxt[ i ];
              break:
            if ( not u_f ) u->nxt[ i ]->fail = rt;
           bfs.push_back( u->nxt[ i ] );
         }
       }
     void match( const string& s, vector< int >& ret ) {
       node* u = rt;
       for ( auto c : s ) {
         while ( u != rt and not u->nxt[ Idx( c ) ] )
           u = u->fail;
          u = u \rightarrow nxt[Idx(c)];
         if ( not u ) u = rt;
         node* tmp = u;
          while ( tmp != rt ) {
            for ( auto d : tmp->data )
              ret.push_back( d );
            tmp = tmp->fail;
       }
} ac;
```

#### 7.4 Suffix Automaton

```
struct Node{
  Node *green, *edge[26];
  int max_len;
  Node(const int _max_len)
    : green(NULL), max_len(_max_len){
    memset(edge,0,sizeof(edge));
} *ROOT, *LAST;
void Extend(const int c) {
  Node *cursor = LAST;
  LAST = new Node((LAST->max_len) + 1);
  for(;cursor&&!cursor->edge[c]; cursor=cursor->green)
    cursor->edge[c] = LAST;
  if (!cursor)
    LAST->green = ROOT;
  else {
    Node *potential_green = cursor->edge[c];
    if((potential_green->max_len)==(cursor->max_len+1))
      LAST->green = potential_green;
    else {
//assert(potential_green->max_len>(cursor->max_len+1));
      Node *wish = new Node((cursor->max_len) + 1);
      for(;cursor && cursor->edge[c]==potential_green;
           cursor = cursor->green)
        cursor->edge[c] = wish;
      for (int i = 0; i < 26; i++)
        wish->edge[i] = potential_green->edge[i];
      wish->green = potential_green->green;
      potential_green->green = wish;
      LAST->green = wish;
```

```
}
char S[10000001], A[10000001];
int N:
int main(){
  scanf("%d%s", &N, S);
  ROOT = LAST = new Node(0);
  for (int i = 0; S[i]; i++)
    Extend(S[i] - 'a');
  while (N--){
  scanf("%s", A);
    Node *cursor = ROOT;
    bool ans = true;
    for (int i = 0; A[i]; i++){
      cursor = cursor->edge[A[i] - 'a'];
      if (!cursor) {
        ans = false;
        break;
    }
    puts(ans ? "Yes" : "No");
  return 0;
```

## 7.5 KMP

```
vector<int> kmp(const string &s) {
  vector<int> f(s.size(), 0);
  /* f[i] = length of the longest prefix
     (excluding s[0:i]) such that it coincides
     with the suffix of s[0:i] of the same length */
   * i + 1 - f[i] is the length of the
     smallest recurring period of s[0:i] */
  int k = 0;
  for (int i = 1; i < (int)s.size(); ++i) {</pre>
   while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
    if (s[i] == s[k]) ++k;
    f[i] = k;
  }
  return f;
vector<int> search(const string &s, const string &t) {
 // return 0-indexed occurrence of t in s
 vector<int> f = kmp(t), res;
for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
    while(k > 0 && (k == (int)t.size() | | s[i] != t[k])
      k = f[k - 1];
    if (s[i] == t[k]) ++k;
    if (k == (int)t.size())
      res.push_back(i - t.size() + 1);
  return res;
```

#### 7.6 Z value

```
char s[MAXN];
int len,z[MAXN];
void Z_value() {
   int i,j,left,right;
   left=right=0; z[0]=len;
   for(i=1;i<len;i++) {
      j=max(min(z[i-left],right-i),0);
      for(;i+j<len&&s[i+j]==s[j];j++);
      z[i]=j;
      if(i+z[i]>right) {
        right=i+z[i];
        left=i;
      }
   }
}
```

#### 7.7 Manacher

```
int z[maxn];
int manacher(const string& s) {
    string t = ".";
    for(char c:s)) t += c, t += '.';
```

```
int 1 = 0, r = 0, ans = 0;
for (int i = 1; i < t.length(); ++i) {
    z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
    while (i - z[i] >= 0 && i + z[i] < t.length()) {
        if(t[i - z[i]] == t[i + z[i]]) ++z[i];
        else break;
    }
    if (i + z[i] > r) r = i + z[i], l = i;
}
for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
    return ans;
}</pre>
```

# 7.8 Lexicographically Smallest Rotation

```
string mcp(string s){
  int n = s.length();
  s += s;
  int i=0, j=1;
  while (i<n && j<n){
    int k = 0;
    while (k < n && s[i+k] == s[j+k]) k++;
    if (s[i+k] <= s[j+k]) j += k+1;
    else i += k+1;
    if (i == j) j++;
  }
  int ans = i < n ? i : j;
  return s.substr(ans, n);
}</pre>
```

# 7.9 BWT

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a'
  vector<int> v[ SIGMA ];
  void BWT(char* ori, char* res){
    // make ori -> ori + ori
    // then build suffix array
  void iBWT(char* ori, char* res){
    for( int i = 0 ; i < SIGMA ; i ++ )</pre>
      v[ i ].clear();
     int len = strlen( ori );
     for( int i = 0 ; i < len ; i ++ )</pre>
      v[ ori[i] - BASE ].push_back( i );
     vector<int> a;
    for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )
for( auto j : v[ i ] ){</pre>
         a.push_back( j );
         ori[ ptr ++ ] = BASE + i;
    for( int i = 0 , ptr = 0 ; i < len ; i ++ ){</pre>
      res[ i ] = ori[ a[ ptr ] ];
      ptr = a[ ptr ];
    res[ len ] = 0;
} bwt;
```

#### 7.10 Palindromic Tree

```
struct palindromic_tree{
  struct node{
    int next[26],f,len;
    int cnt,num,st,ed;
    node(int l=0):f(0),len(1),cnt(0),num(0){
      memset(next, 0, sizeof(next));
    }
  };
  vector<node> state;
  vector<char> s;
  int last,n;
  void init(){
    state.clear();s.clear();last=1; n=0;
    state.push_back(0);state.push_back(-1);
    state[0].f=1;s.push_back(-1);
  int getFail(int x){
    while(s[n-state[x].len-1]!=s[n])x=state[x].f;
```

```
return x:
  void add(int c){
    s.push_back(c-='a'); ++n;
    int cur=getFail(last);
    if(!state[cur].next[c]){
      int now=state.size();
      state.push_back(state[cur].len+2);
      state[now].f=state[getFail(state[cur].f)].next[c
          ];
      state[cur].next[c]=now;
      state[now].num=state[state[now].f].num+1;
    last=state[cur].next[c];
    ++state[last].cnt;
  int size(){
    return state.size()-2;
} pt;
int main() {
 string s; cin >> s; pt.init();
 for (int i=0; i<SZ(s); i++) {</pre>
    int prvsz = pt.size(); pt.add(s[i]);
    if (prvsz != pt.size()) {
      int r = i, l = r - pt.state[pt.last].len + 1;
      // pal @ [l,r]: s.substr(l, r-l+1)
    }
  return 0;
```

# 8 Misc

#### 8.1 Theorems

#### 8.1.1 Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii}=d(i)$ ,  $L_{ij}=-c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|{\rm det}(\tilde{L}_{rr})|$  .

#### 8.1.2 Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

#### 8.1.3 Cayley's Formula

- Given a degree sequence  $d_1,d_2,\ldots,d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .

#### 8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all  $1 \leq k \leq n$ .

# 8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

#### 8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let  $N_G(W)$  denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff  $\forall W\subseteq X, |W|\leq |N_G(W)|$ 

# 8.1.7 Euler's planar graph formula

```
V - E + F = C + 1, E \le 3V - 6(?)
```

#### 8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have  $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$ 

### 8.2 MaximumEmptyRect

```
int max_empty_rect(int n, int m, bool blocked[N][N]){
  static int mxu[2][N], me=0,he=1,ans=0;
  for(int i=0;i<m;i++) mxu[he][i]=0;</pre>
  for(int i=0;i<n;i++){</pre>
    stack<PII,vector<PII>> stk;
    for(int j=0;j<m;++j){</pre>
      if(blocked[i][j]) mxu[me][j]=0;
      else mxu[me][j]=mxu[he][j]+1;
      int la = j;
      while(!stk.empty()&&stk.top().FF>mxu[me][j]){
        int x1 = i - stk.top().FF, x2 = i;
        int y1 = stk.top().SS, y2 = j;
        la = stk.top().SS; stk.pop();
        ans=\max(ans,(x2-x1)*(y2-y1));
      if(stk.empty()||stk.top().FF<mxu[me][j])</pre>
        stk.push({mxu[me][j],la});
    while(!stk.empty()){
      int x1 = i - stk.top().FF, x2 = i;
      int y1 = stk.top().SS-1, y2 = m-1;
      stk.pop();
      ans=max(ans,(x2-x1)*(y2-y1));
    swap(me,he);
  return ans;
}
```

# 8.3 DP-opt Condition

## 8.3.1 totally monotone (concave/convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j' \text{, } B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

# 8.3.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j' \text{, } B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

## 8.4 Convex 1D/1D DP

```
struct segment {
  int i, 1, r;
  segment() {}
  segment(int a, int b, int c): i(a), l(b), r(c) {}
inline lld f(int l, int r) {return dp[l] + w(l + 1, r)
    ;}
void solve() {
  dp[0] = 011;
  deque<segment> deq; deq.push_back(segment(0, 1, n));
  for (int i = 1; i <= n; ++i) {</pre>
    dp[i] = f(deq.front().i, i);
    while(deq.size()&&deq.front().r<i+1)deq.pop_front()</pre>
    deq.front().l = i + 1;
    segment seg = segment(i, i + 1, n);
    while (deq.size() &&
       f(i, deq.back().1) < f(deq.back().i, deq.back().1))
         deq.pop_back();
    if (deq.size()) {
      int d = 1048576, c = deq.back().1;
while (d >>= 1) if (c + d <= deq.back().r)</pre>
         if(f(i, c + d) > f(deq.back().i, c + d)) c += d
      deq.back().r = c; seg.l = c + 1;
    if (seg.l <= n) deq.push_back(seg);</pre>
}
```

### 8.5 ConvexHull Optimization

```
inline 1ld DivCeil(1ld n, 1ld d) { // ceil(n/d)
  return n / d + (((n < 0) != (d > 0)) && (n % d));
struct Line {
  static bool flag;
  11d a, b, 1, r; // y=ax+b in [l, r)
  11d operator()(11d x) const { return a * x + b; }
  bool operator<(const Line& i) const {</pre>
    return flag ? tie(a, b) < tie(i.a, i.b) : 1 < i.l;</pre>
  11d operator&(const Line& i) const {
    return DivCeil(b - i.b, i.a - a);
  }
bool Line::flag = true;
class ConvexHullMax {
  set<Line> L;
 public:
  ConvexHullMax() { Line::flag = true; }
  void InsertLine(lld a, lld b) { // add y = ax + b
    Line now = {a, b, -INF, INF};
    if (L.empty()) {
      L.insert(now);
      return;
    Line::flag = true;
    auto it = L.lower_bound(now);
    auto prv = it == L.begin() ? it : prev(it);
    if (it != L.end() && ((it != L.begin() &&
       (*it)(it->1) >= now(it->1) &&
       (*prv)(prv->r - 1) >= now(prv->r - 1)) ||
       (it == L.begin() && it->a == now.a))) return;
    if (it != L.begin()) {
      while (prv != L.begin() &&
         (*prv)(prv->1) <= now(prv->1))
          prv = --L.erase(prv);
      if (prv == L.begin() && now.a == prv->a)
        L.erase(prv);
    if (it != L.end())
       while (it != --L.end() &&
        (*it)(it->r) \leftarrow now(it->r)
          it = L.erase(it);
    if (it != L.begin()) {
      prv = prev(it);
       const_cast<Line*>(&*prv)->r=now.l=((*prv)&now);
    if (it != L.end())
      const_cast<Line*>(&*it)->l=now.r=((*it)&now);
    L.insert(it, now);
  11d Query(11d a) const { // query max at x=a
    if (L.empty()) return -INF;
    Line::flag = false;
    auto it = --L.upper_bound({0, 0, a, 0});
    return (*it)(a);
|};
```

# 8.6 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
   int s = 0;
   for (int i = 2; i <= n; i++)
        s = (s + m) % i;
   return s;
}
// died at kth
int kth(int n, int m, int k){
   if (m == 1) return n-1;
   for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
   return k;
}
```

# 8.7 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
```

```
void tarjan(int u){
  dfn[u]=low[u]=++dfs_idx;
  for(int i=0;i<(int)init_g[u].size();i++){</pre>
    int v=init_g[u][i];
    if(v==par[u]) continue;
    if(!dfn[v]){
      par[v]=u;
      tarjan(v);
      low[u]=min(low[u],low[v]);
      if(dfn[u]<low[v]){</pre>
        g[u].push_back(v);
        g[v].push_back(u);
    }else{
      low[u]=min(low[u],dfn[v]);
      if(dfn[v]<dfn[u]){</pre>
        int temp_v=u;
        bcc id++;
        while(temp_v!=v){
          g[bcc_id+n].push_back(temp_v);
          g[temp_v].push_back(bcc_id+n);
          temp_v=par[temp_v];
        g[bcc_id+n].push_back(v);
        g[v].push_back(bcc_id+n);
        reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
    }
 }
int dp[maxn][2],min_dp[2][2],tmp[2][2],tp[2];
void dfs(int u,int fa){
 if(u<=n){
    for(int i=0;i<(int)g[u].size();i++){</pre>
      int v=g[u][i];
      if(v==fa) continue;
      dfs(v,u);
      memset(tp,0x8f,sizeof tp);
      if(v<=n){
        tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
        tp[1]=max(
          dp[u][0]+dp[v][0]+1,
          dp[u][1]+max(dp[v][0],dp[v][1])
        );
      }else{
        tp[0]=dp[u][0]+dp[v][0];
        tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
      dp[u][0]=tp[0],dp[u][1]=tp[1];
  }else{
    for(int i=0;i<(int)g[u].size();i++){</pre>
      int v=g[u][i];
      if(v==fa) continue;
      dfs(v,u);
    min dp[0][0]=0;
    min_dp[1][1]=1;
    for(int i=0;i<(int)g[u].size();i++){</pre>
      int v=g[u][i];
      if(v==fa) continue;
      memset(tmp,0x8f,sizeof tmp);
      tmp[0][0]=max(
        min_dp[0][0]+max(dp[v][0],dp[v][1]),
        min_dp[0][1]+dp[v][0]
      );
      tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
      tmp[1][0]=max(
        min_dp[1][0]+max(dp[v][0],dp[v][1]),
        min_dp[1][1]+dp[v][0]
      tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
      memcpy(min_dp,tmp,sizeof tmp);
    dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
    dp[u][0]=min_dp[0][0];
  }
int main(){
  int m,a,b;
  scanf("%d%d",&n,&m);
```

resume(c);

```
for(int i=0;i<m;i++){</pre>
                                                             } sol;
    scanf("%d%d",&a,&b);
    init_g[a].push_back(b);
                                                              8.9 Tree Knapsack
    init_g[b].push_back(a);
                                                              int dp[N][K];PII obj[N];
 par[1]=-1;
                                                              vector<int> G[N];
  tarjan(1);
                                                              void dfs(int u, int mx){
  dfs(1,-1);
                                                                 for(int s: G[u]) {
 printf("%d\n", max(dp[1][0], dp[1][1]));
                                                                   if(mx < obj[s].first) continue;</pre>
  return 0;
                                                                   for(int i=0;i<=mx-obj[s].FF;i++)</pre>
                                                                   dp[s][i] = dp[u][i];
dfs(s, mx - obj[s].first);
                                                                   for(int i=obj[s].FF;i<=mx;i++)</pre>
8.8 DLX
                                                                     dp[u][i] = max(dp[u][i],
                                                                       dp[s][i - obj[s].FF] + obj[s].SS);
struct DLX {
                                                                 }
  const static int maxn=210;
                                                              }
  const static int maxm=210;
                                                              int main(){
  const static int maxnode=210*210;
                                                                int n, k; cin >> n >> k;
 int n, m, size, row[maxnode], col[maxnode];
int U[maxnode], D[maxnode], L[maxnode], R[maxnode];
                                                                 for(int i=1;i<=n;i++){</pre>
                                                                   int p; cin >> p;
  int H[maxn], S[maxm], ansd, ans[maxn];
                                                                   G[p].push_back(i);
  void init(int _n, int _m) {
                                                                   cin >> obj[i].FF >> obj[i].SS;
   n = _n, m = _m;
for(int i = 0; i <= m; ++i) {</pre>
                                                                 dfs(0, k); int ans = 0;
      S[i] = 0;
                                                                 for(int i=0;i<=k;i++) ans = max(ans, dp[0][i]);</pre>
      U[i] = D[i] = i;
                                                                 cout << ans << '\n';
      L[i] = i-1, R[i] = i+1;
                                                                 return 0:
    R[L[0] = size = m] = 0;
    for(int i = 1; i <= n; ++i) H[i] = -1;</pre>
                                                              8.10 N Queens Problem
  void Link(int r, int c) {
    ++S[col[++size] = c];
                                                              vector< int > solve( int n ) {
    row[size] = r; D[size] = D[c];
                                                                 // no solution when n=2, 3
    U[D[c]] = size; U[size] = c; D[c] = size;
                                                                 vector< int > ret;
    if(H[r] < 0) H[r] = L[size] = R[size] = size;</pre>
                                                                 if ( n % 6 == 2 ) {
    else {
                                                                   for ( int i = 2 ; i <= n ; i += 2 )</pre>
      R[size] = R[H[r]];
                                                                     ret.push_back( i );
      L[R[H[r]]] = size;
                                                                   ret.push_back( 3 ); ret.push_back( 1 );
      L[size] = H[r];
                                                                   for ( int i = 7 ; i <= n ; i += 2 )</pre>
      R[H[r]] = size;
                                                                     ret.push_back( i );
    }
                                                                   ret.push_back( 5 );
                                                                 } else if ( n % 6 == 3 ) {
  void remove(int c) {
                                                                   for ( int i = 4 ; i <= n ; i += 2 )
    L[R[c]] = L[c]; R[L[c]] = R[c];
                                                                     ret.push_back( i );
    for(int i = D[c]; i != c; i = D[i])
                                                                   ret.push_back( 2 );
      for(int j = R[i]; j != i; j = R[j]) {
                                                                   for ( int i = 5 ; i <= n ; i += 2 )</pre>
        U[D[j]] = U[j];
                                                                     ret.push_back( i );
        D[U[j]] = D[j];
                                                                   ret.push_back( 1 ); ret.push_back( 3 );
        --S[col[j]];
                                                                 } else {
      }
                                                                   for ( int i = 2 ; i <= n ; i += 2 )
                                                                     ret.push_back( i );
  void resume(int c) {
                                                                   for ( int i = 1 ; i <= n ; i += 2 )</pre>
    L[R[c]] = c; R[L[c]] = c;
                                                                     ret.push_back( i );
    for(int i = U[c]; i != c; i = U[i])
      for(int j = L[i]; j != i; j = L[j]) {
                                                                 return ret;
        U[D[j]] = j;
                                                              }
        D[U[j]] = j;
        ++S[col[j]];
    }
  void dance(int d) {
    if(d>=ansd) return;
    if(R[0] == 0) {
  ansd = d;
      return;
    int c = R[0];
    for(int i = R[0]; i; i = R[i])
      if(S[i] < S[c]) c = i;
    remove(c);
    for(int i = D[c]; i != c; i = D[i]) {
      ans[d] = row[i];
      for(int j = R[i]; j != i; j = R[j])
        remove(col[j]);
      dance(d+1);
      for(int j = L[i]; j != i; j = L[j])
        resume(col[j]);
```