Contents

1	Basi																						1
	1.1	vimrc																					1
	1.2	Increase Stack																					1
	1.3	Pragma optimiz	ation																				1
	1.4	IO Optimization																					1
2	Date	3 Structure																					2
	2.1	Bigint																					2
	2.2	Dark Magic																					2
	2.3	Disjoint Set																					3
	2.4	Link-Cut Tree .																					3
	2.5	LiChao Segmen																					4
		-																					4
	2.6	Treap																					
	2.7	SparseTable																					4
	2.8	Linear Basis			٠	٠		•		•	٠		 ٠	•	•	٠						 ٠	4
_	_																						_
3	Grap																						5
	3.1	Euler Circuit																					5
	3.2	BCC Edge																					5
	3.3	BCC Vertex																					5
	3.4	2-SAT (SCC)																					5
	3.5	Lowbit Decomp																					6
	3.6	MaxClique																					6
	3.7	Virtural Tree .																					7
	3.8																						7
	3.9	Tree Hashing . Minimum Mean																					7
			_																				
		Mo's Algorithm																					7
	3.11	Minimum Steine																					8
		Directed Minimu																					8
	3.13	Dominator Tree																					9
4	Mate	ching & Flow																					9
	4.1	Kuhn Munkres .																					9
	4.2	Bipartite Matchi	ng																				9
	4.3	General Graph I	Vatchir	ıq.																			10
	4.4	Minimum Weigh																					10
	4.5	Flow Models																					10
	4.6	Dinic																					11
	4.7	Minimum Cost N																					11
	4.8	Global Min-Cut																					12
	4.8	Global Min-Cut			•	٠		•	•	•	•	•	 •	•	•	•	•	•	•	•	•	 •	12
5	Matl																						12
5																							
	5.1	Prime Table																					12
	5.2	$\lfloor \frac{n}{i} \rfloor$ Enumeration																					12
	5.3	5 5																					12
	5.4	Pollard Rho																					12
	5.5	Pi Count (Linear	Sieve)																				12
	E (12
	5.6	Range Sieve																					
	5.7	Range Sieve Miller Rabin																					13
	5.7	Miller Rabin																					13
	5.7 5.8	Miller Rabin Inverse Element											 :										13 13
	5.7 5.8 5.9	Miller Rabin Inverse Element Euler Phi Function	 : on				 						 										13 13 13
	5.7 5.8 5.9 5.10	Miller Rabin Inverse Element Euler Phi Functio Gauss Eliminatio	 : on				 						 										13 13 13 13
	5.7 5.8 5.9 5.10 5.11	Miller Rabin Inverse Element Euler Phi Functio Gauss Eliminatio Fast Fourier Tra	on on on	 			· · · · · · · ·						 									 	13 13 13 13 13
	5.7 5.8 5.9 5.10 5.11 5.12	Miller Rabin Inverse Element Euler Phi Functio Gauss Eliminatio Fast Fourier Tra High Speed Line	on on on	 urr	en								 									 	13 13 13 13 13 13
	5.7 5.8 5.9 5.10 5.11 5.12 5.13	Miller Rabin Inverse Element Euler Phi Functio Gauss Eliminatio Fast Fourier Tra High Speed Line Chinese Remain	on	 urr	en								 									 	13 13 13 13 13 13 14
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Tra High Speed Line Chinese Remain Berlekamp Mas	on on on	 urr	en								 									 	13 13 13 13 13 13 14 14
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Tra High Speed Line Chinese Remain Berlekamp Mas.	on		en																		13 13 13 13 13 13 14 14 14
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Tra High Speed Line Chinese Remain Berlekamp Mas	on		en																		13 13 13 13 13 13 14 14
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Tra High Speed Line Chinese Remain Berlekamp Mas.	on		en								 									 	13 13 13 13 13 13 14 14 14
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Tra High Speed Line Chinese Remain Berlekamp Mass NTT Polynomial Ope	on		en																	 	13 13 13 13 13 14 14 14 15
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Tradingh Speed Line Chinese Remain Berlekamp Massint Polynomial Ope FWT DiscreteLog	on																			 	13 13 13 13 13 14 14 14 15 16
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas. NTT Polynomial Ope FWT	on		en																		13 13 13 13 13 14 14 14 15 16
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Tra High Speed Line Chinese Remain Berlekamp Mas: NTT Polynomial Ope FWT	on		en																		13 13 13 13 13 14 14 15 16 16 16
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatio Fast Fourier Tra High Speed Line Chinese Remain Berlekamp Mas. NTT Polynomial Ope FWT	on		en																		13 13 13 13 13 14 14 15 16 16 16 16
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Tra High Speed Line Chinese Remain Berlekamp Mas: NTT Polynomial Ope FWT	on		en																		13 13 13 13 13 14 14 15 16 16 16
6	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mass NTT Polynomial Ope FWT DiscreteLog Quadratic residu De-Bruijn Simplex Construs Simplex	on		en																		13 13 13 13 13 14 14 15 16 16 16 16 17
6	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22 Geol	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas. NTT Polynomial Ope FWT	on		en																		13 13 13 13 13 14 14 15 16 16 16 16 17
6	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 5.21 5.22 Geol	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Tra High Speed Line Chinese Remain Berlekamp Mas: NTT Polynomial Ope FWT	on		en																		13 13 13 13 13 14 14 15 16 16 16 16 17 17
6	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 5.21 5.22 Geol	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas. NTT	on		en																		13 13 13 13 13 14 14 15 16 16 16 16 17 17
6	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 5.21 5.22 Geol	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas: NTT Polynomial Ope FWT	on		en																		13 13 13 13 13 14 14 14 15 16 16 16 17 17 17 17 18
6	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.20 5.21 5.22 Geol 6.1 6.2 6.3 6.4	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mass NTT Polynomial Ope FWT DiscreteLog Quadratic reside De-Bruijn Simplex Constructions Simplex	on		en																		13 13 13 13 13 14 14 14 15 16 16 16 17 17 17 17 18 18
6	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 5.21 5.22 Geol 6.2 6.3 6.4 6.5	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas NTT	on		en																		13 13 13 13 13 13 14 14 14 15 16 16 16 16 17 17 17 17 18 18 18
6	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 5.21 5.22 Geo 6.1 6.2 6.3 6.4 6.5 6.6	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas. NTT	on		en																		13 13 13 13 13 13 14 14 14 15 16 16 16 16 17 17 17 17 18 18 18 18
6	5.7 5.8 5.9 5.10 5.11 5.12 5.14 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas. NTT	on		en																		13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17 17 17 18 18 18 18 19
6	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mass. NTT Polynomial Ope FWT	on		en																		13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17 17 17 18 18 18 18 19 19
6	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mass NTT Polynomial Ope FWT	on		en																		13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17 17 17 18 18 18 18 19
6	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas NTT	on																				13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17 17 17 18 18 18 18 19 19
6	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas NTT	on																				13 13 13 13 13 14 14 14 15 16 16 16 17 17 17 17 18 18 18 19 19
6	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas. NTT	on		en																		13 13 13 13 13 14 14 14 15 16 16 16 17 17 17 17 17 18 18 18 18 19 19 19
6	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.16 5.17 5.18 5.19 5.21 5.22 Geo 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 6.11 6.12	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas. NTT	on		en r .																		13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17 17 18 18 18 18 19 19 19 19 19 19 19 19 19 19 19 19 19
6	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.16 5.17 5.18 5.19 5.21 5.22 Geo 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 6.11 6.12	Miller Rabin Inverse Element Euler Phi Functic Gauss Eliminatic Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mas. NTT	on		en r .																		13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17 17 18 18 18 18 19 19 19 19 19 19 19 19 19 19 19 19 19
66	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 6.12 6.13	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mass NTT Polynomial Ope FWT	on		en r .																		13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17 17 17 18 18 18 18 19 19 19 19 19 19 19 19 19 19 19 19 19
	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.10 6.11 6.12 6.13 5.13 6.14 6.15 6.15 6.16 6.17 6.17 6.17 6.17 6.17 6.17 6.17	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mass NTT Polynomial Ope FWT	on	ge																			13 13 13 13 13 13 14 14 14 15 16 16 16 16 17 17 17 17 17 18 18 18 18 19 19 19 19 19 19 19 19 19 19 19 19 19
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 6.12 6.13 6.14 6.15 6.15 6.16 6.17 6.17 6.17 6.18 6.19 6.19 6.19 6.19 6.19 6.19 6.19 6.19	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trailing Speed Line Chinese Remain Berlekamp Mass. NTT	nsformer Recorder																				13 13 13 13 13 13 14 14 15 16 16 16 16 16 17 17 17 17 18 18 18 19 19 19 19 20 20 20 20 20 20 20 20 20 20 20 20 20
	5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 6.12 6.13 7.1 7.1 7.1	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Track High Speed Line Chinese Remain Berlekamp Mas. NTT	on																				13 13 13 13 13 13 14 14 15 16 16 16 16 16 17 17 17 17 17 18 18 18 18 19 19 19 20 20 20 20 20 20 20 20 20 20 20 20 20
	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.10 6.12 6.13 Strir 7.1 7.2 7.3	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mass NTT Polynomial Ope FWT	nsform car Recoder																				13 13 13 13 13 14 14 14 15 16 16 16 17 17 17 17 17 18 18 18 18 19 19 19 19 20 20 20 20 20 20 20 20 20 20 20 20 20
	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.10 6.11 6.13 Strir 7.1 7.2 7.3 7.4	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mass NTT Polynomial Ope FWT	on																				13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17 17 17 18 18 18 19 19 19 20 20 20 20 21 21 21 21 21 21 21 21 21 21 21 21 21
	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 6.3 6.4 6.5 6.6 6.7 6.8 6.10 6.11 6.12 6.13 7.1 7.2 7.3 7.4 7.5	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mass NTT	on	ge																			13 13 13 13 13 14 14 14 15 16 16 16 16 17 17 17 17 18 18 18 19 19 19 20 20 20 20 21 21 22 22 22 22 22 22 22 22 22 22 22
	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.21 5.22 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 6.12 6.13 7.1 7.2 7.3 7.7 7.7 7.7 7.7 7.7 7.7 7.7 7.7 7.7	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trailing Speed Line Chinese Remain Berlekamp Mas. NTT	rations centre control	geele																			13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17 17 18 18 18 19 19 19 20 20 20 21 22 22 22 22 22 22 22 22 22 22 22 22
	5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 6.3 6.4 6.5 6.6 6.7 6.8 6.10 6.11 6.12 6.13 7.1 7.2 7.3 7.4 7.5	Miller Rabin Inverse Element Euler Phi Function Gauss Elimination Fast Fourier Trad High Speed Line Chinese Remain Berlekamp Mass NTT	nsform ar Recoder sey so sey so section sectio	on the second se																			13 13 13 13 13 14 14 14 15 16 16 16 16 17 17 17 17 18 18 18 19 19 19 20 20 20 20 21 21 22 22 22 22 22 22 22 22 22 22 22

```
23
8 Misc
 23
  8.1.4 Erdős-Gallai theorem
8.1.5 Havel-Hakimi algorithm
8.1.6 Hall's marriage theorem
  8.1.6 Hairs marriage theorem
8.1.7 Euler's planar graph formula
8.1.8 Pick's theorem
8.1.9 Lucas's theorem
 8.3.1 totally monotone (concave/convex) . . . . . . . . . . . . . . .
  8.3.2 monge condition (concave/convex) . . . . . . . . . . . . . . .
 1
  Basic
```

1.1 vimrc

```
se is nu rnu bs=2 ru mouse=a encoding=utf-8
se cin et ts=4 sw=4 sts=4 t_Co=256
syn on
colorscheme ron
filetype indent on
map <F8> <ESC>:w<CR>:!clear && g++ "%" -o "%<" -
    fsanitize=address -fsanitize=undefined -g && echo
    success<CR>
map <F9> <ESC>:w<CR>:!clear && g++ "%" -o "%<" -02 &&
    echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
```

1.2 Increase Stack

```
const int size = 256 << 20;</pre>
register long rsp asm("rsp");
char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));
```

1.3 Pragma optimization

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,sse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
```

1.4 IO Optimization

```
static inline int gc() {
  static char buf[ 1 << 20 ], *p = buf, *end = buf;</pre>
  if ( p == end ) {
    end = buf + fread( buf, 1, 1 << 20, stdin );
     if ( end == buf ) return EOF;
    p = buf;
  return *p++;
template < typename T >
static inline bool gn( T &_ ) {
  register int c = gc(); register T __ = 1; _ = 0;
while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
if(c == '-') { __ = -1; c = gc(); }
  if(c == EOF) return false;
  while('0'<=c&&c<='9') _{-} = _{-} * 10 + c _{-} '0', c = gc();
  _ *= __;
  return true;
template < typename T, typename ...Args >
static inline bool gn( T &x, Args &...args )
{ return gn(x) && gn(args...); }
```

2 Data Structure

2.1 Bigint

```
class BigInt{
private:
 using 1ld = int_fast64_t;
 #define PRINTF_ARG PRIdFAST64
 #define LOG_BASE_STR "9
  static constexpr lld BASE = 1000000000;
 static constexpr int LOG_BASE = 9;
 vector<lld> dig; bool neg;
 inline int len() const { return (int) dig.size(); }
 inline int cmp_minus(const BigInt& a) const {
    if(len() == 0 && a.len() == 0) return 0;
    if(neg ^ a.neg)return a.neg ^ 1;
    if(len()!=a.len())
      return neg?a.len()-len():len()-a.len();
    for(int i=len()-1;i>=0;i--) if(dig[i]!=a.dig[i])
      return neg?a.dig[i]-dig[i]:dig[i]-a.dig[i];
    return 0;
 inline void trim(){
    while(!dig.empty()&&!dig.back())dig.pop_back();
    if(dig.empty()) neg = false;
public:
 BigInt(): dig(vector<lld>()), neg(false){}
 BigInt(lld a): dig(vector<lld>()){
   neg = a<0; dig.push_back(abs(a));</pre>
    trim();
 BigInt(const string& a): dig(vector<lld>()){
    assert(!a.empty()); neg = (a[0]=='-');
    for(int i=((int)a.size())-1;i>=neg;i-=LOG_BASE){
      11d cur = 0;
      for(int j=min(LOG_BASE-1,i-neg);j>=0;j--)
        cur = cur*10+a[i-j]-'0';
      dig.push_back(cur);
    } trim();
  inline bool operator<(const BigInt& a)const
    {return cmp_minus(a)<0;}
  inline bool operator<=(const BigInt& a)const
    {return cmp_minus(a)<=0;}
  inline bool operator==(const BigInt& a)const
    {return cmp_minus(a)==0;}
  inline bool operator!=(const BigInt& a)const
    {return cmp_minus(a)!=0;}
  inline bool operator>(const BigInt& a)const
    {return cmp_minus(a)>0;}
  inline bool operator>=(const BigInt& a)const
    {return cmp_minus(a)>=0;}
 BigInt operator-() const {
    BigInt ret = *this;
    ret.neg ^= 1; return ret;
 BigInt operator+(const BigInt& a) const {
    if(neg) return -(-(*this)+(-a));
    if(a.neg) return (*this)-(-a);
    int n = max(a.len(), len())
    BigInt ret; ret.dig.resize(n);
    11d pro = 0;
    for(int i=0;i<n;i++) {</pre>
      ret.dig[i] = pro;
      if(i < a.len()) ret.dig[i] += a.dig[i];</pre>
      if(i < len()) ret.dig[i] += dig[i];</pre>
      pro = 0:
      if(ret.dig[i] >= BASE) pro = ret.dig[i]/BASE;
      ret.dig[i] -= BASE*pro;
    if(pro != 0) ret.dig.push_back(pro);
    return ret;
 BigInt operator-(const BigInt& a) const {
    if(neg) return -(-(*this) - (-a));
    if(a.neg) return (*this) + (-a);
    int diff = cmp_minus(a);
    if(diff < 0) return -(a - (*this));</pre>
    if(diff == 0) return 0;
    BigInt ret; ret.dig.resize(len(), 0);
    for(int i=0;i<len();i++) {</pre>
```

```
ret.dig[i] += dig[i];
      if(i < a.len()) ret.dig[i] -= a.dig[i];</pre>
      if(ret.dig[i] < 0){</pre>
        ret.dig[i] += BASE;
        ret.dig[i+1]--;
    ret.trim(); return ret;
  BigInt operator*(const BigInt& a) const {
    if(!len()||!a.len()) return 0;
    BigInt ret; ret.dig.resize(len()+a.len()+1);
    ret.neg = neg ^ a.neg;
    for(int i=0;i<len();i++)</pre>
      for(int j=0;j<a.len();j++){</pre>
        ret.dig[i+j] += dig[i] * a.dig[j];
        if(ret.dig[i+j] >= BASE) {
          lld x = ret.dig[i+j] / BASE;
          ret.dig[i+j+1] += x;
          ret.dig[i+j] -= x * BASE;
        }
    ret.trim(); return ret;
  BigInt operator/(const BigInt& a) const {
    assert(a.len());
    if(len() < a.len()) return 0;</pre>
    BigInt ret; ret.dig.resize(len()-a.len()+1);
    ret.neg = a.neg;
    for(int i=len()-a.len();i>=0;i--){
      11d 1 = 0, r = BASE;
      while(r-l > 1){
        1ld mid_=_(l+r)>>1;
        ret.dig[i] = mid;
        if(ret*a<=(neg?-(*this):(*this))) 1 = mid;</pre>
        else r = mid;
      ret.dig[i] = 1;
    ret.neg ^= neg; ret.trim();
    return ret;
  BigInt operator%(const BigInt& a) const {
    return (*this) - (*this) / a * a;
  friend BigInt abs(BigInt a) { a.neg = 0; return a; }
  friend void swap(BigInt& a, BigInt& b){
    swap(a.dig, b.dig); swap(a.neg, b.neg);
  friend istream& operator>>(istream& ss, BigInt& a){
    string s; ss >> s; a = s; return ss;
  friend ostream&operator<<(ostream&o, const BigInt&a){</pre>
    if(a.len() == 0) return o << '0';
if(a.neg) o << '-';
    ss << o.dig.back();
    for(int i=a.len()-2;i>=0;i--)
      o<<setw(LOG_BASE)<<setfill('0')<<a.dig[i];
    return o;
  inline void print() const {
    if(len() == 0){putchar('0');return;}
    if(neg) putchar('-');
printf("%" PRINTF_ARG, dig.back());
    for(int i=len()-2;i>=0;i--)
      printf("%0" LOG_BASE_STR PRINTF_ARG, dig[i]);
  #undef PRINTF_ARG
  #undef LOG_BASE_STR
2.2 Dark Magic
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using __gnu_pbds::pairing_heap_tag;
using __gnu_pbds::binary_heap_tag;
using __gnu_pbds::binomial_heap_tag;
```

using __gnu_pbds::rc_binomial_heap_tag;
using __gnu_pbds::thin_heap_tag;

template<typename T>

inline bool is_root(){

```
using pbds_heap=__gnu_pbds::prioity_queue<T, less<T>, \
                                                                 return par==nullptr ||\
                                                                   (par->ch[0]!=this && par->ch[1]!=this);
                                       pairing_heap_tag>;
// a.join(b), pq.modify(pq.push(10),
using __gnu_pbds::rb_tree_tag;
                                                               bool is_rch(){return !is_root() && par->ch[1]==this;}
using __gnu_pbds::ov_tree_tag;
                                                             } *node[maxn],*stk[maxn];
using __gnu_pbds::splay_tree_tag;
                                                             int top;
                                                             void to_child(Node* p,Node* c,bool dir){
template<typename T>
using ordered_set = __gnu_pbds::tree<T,\</pre>
                                                               p->ch[dir]=c;
__gnu_pbds::null_type,less<T>,rb_tree_tag,\
                                                               p->up();
__gnu_pbds::tree_order_statistics_node_update>;
                                                             inline void rotate(Node* node){
// find_by_order, order_of_key
template<typename A, typename B>
                                                               Node* par=node->par;
using hTable1=__gnu_pbds::cc_hash_table<A,B>;
                                                               Node* par_par=par->par;
template<typename A, typename B>
                                                               bool dir=node->is_rch();
using hTable2=__gnu_pbds::gp_hash_table<A,B>;
                                                               bool par_dir=par->is_rch()
                                                               to_child(par, node->ch[!dir], dir);
                                                               to_child(node,par,!dir);
2.3
      Disjoint Set
                                                               if(par_par!=nullptr && par_par->ch[par_dir]==par)
                                                                 to_child(par_par, node, par_dir);
class DJS {
                                                               else node->par=par_par;
private:
  vector< int > fa, sz, sv;
                                                             inline void splay(Node* node){
  vector< pair< int*, int > > opt;
                                                               Node* tmp=node;
  void assign( int *k, int v ) {
                                                               stk[top++]=node;
    opt.emplace_back( k, *k );
                                                               while(!tmp->is_root()){
    *k = v:
                                                                 tmp=tmp->par;
                                                                 stk[top++]=tmp;
public:
  void init( int n ) {
                                                               while(top) stk[--top]->down();
    fa.resize( n ); iota( fa.begin(), fa.end(), 0 );
sz.resize( n ); fill( sz.begin(), sz.end(), 1 );
                                                               for(Node *fa=node->par;
                                                                !node->is_root();
    opt.clear();
                                                                rotate(node), fa=node->par)
                                                                 if(!fa->is_root())
  int query(int x) {return fa[x] == x?x:query(fa[x]);}
                                                                   rotate(fa->is_rch()==node->is_rch()?fa:node);
  void merge( int a, int b ) {
    int af = query( a ), bf = query( b );
                                                             inline void access(Node* node){
    if( af == bf ) return;
                                                               Node* last=nullptr;
    if( sz[ af ] < sz[ bf ] ) swap( af, bf );</pre>
                                                               while(node!=nullptr){
    assign( &fa[ bf ], fa[ af ] );
assign( &sz[ af ], sz[ af ] + sz[ bf ] );
                                                                 splay(node)
                                                                 to_child(node, last, true);
                                                                 last=node:
  void save() { sv.push_back( (int) opt.size() ); }
                                                                 node=node->par;
  void undo() {
    int ls = sv.back(); sv.pop_back();
    while ( ( int ) opt.size() > ls )
                                                             inline void change_root(Node* node){
      pair< int*, int > cur = opt.back();
                                                               access(node);splay(node);node->set_rev();
      *cur.first = cur.second;
      opt.pop_back();
                                                             inline void link(Node* x,Node* y){
                                                               change_root(x);splay(x);x->par=y;
};
                                                             inline void split(Node* x, Node* y){
                                                               change_root(x);access(y);splay(x);
2.4
      Link-Cut Tree
                                                               to_child(x,nullptr,true);y->par=nullptr;
                                                             inline void change_val(Node* node,int v){
struct Node{
  Node *par, *ch[2];
                                                               access(node);splay(node);node->v=v;node->up();
  int xor_sum, v;
                                                             inline int query(Node* x, Node* y){
  bool is_rev;
                                                               change_root(x);access(y);splay(y);
  Node(int _v){
                                                               return y->xor_sum;
    v=xor_sum=_v;is_rev=false;
    par=ch[0]=ch[1]=nullptr;
                                                             inline Node* find_root(Node* node){
                                                               access(node);splay(node);
  inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
  inline void down(){
                                                               Node* last=nullptr;
    if(is_rev){
                                                               while(node!=nullptr){
      if(ch[0]!=nullptr) ch[0]->set_rev();
                                                                 node->down();last=node;node=node->ch[0];
      if(ch[1]!=nullptr) ch[1]->set_rev();
      is_rev=false;
                                                               return last;
    }
                                                             set<pii> dic;
                                                             inline void add_edge(int u,int v){
  inline void up(){
    xor_sum=v;
                                                               if(u>v) swap(u,v)
    if(ch[0]!=nullptr){
                                                               if(find_root(node[u])==find_root(node[v])) return;
                                                               dic.insert(pii(u,v))
      xor_sum^=ch[0]->xor_sum;
      ch[0]->par=this;
                                                               link(node[u],node[v]);
                                                             inline void del_edge(int u,int v){
    if(ch[1]!=nullptr){
      xor_sum^=ch[1]->xor_sum;
                                                               if(u>v) swap(u,v);
                                                               if(dic.find(pii(u,v))==dic.end()) return;
      ch[1]->par=this;
                                                               dic.erase(pii(u,v))
                                                               split(node[u], node[v]);
```

2.5 LiChao Segment Tree

```
struct Line{
 int m, k, id;
Line() : id( -1 ) {}
  Line( int a, int b, int c )
    : m( a ), k( b ), id( c ) {}
  int at( int x ) { return m * x + k; }
class LiChao {
  private:
    int n; vector< Line > nodes;
    inline int lc( int x ) { return 2 * x + 1; }
    inline int rc( int x ) { return 2 * x + 2;
    void insert( int 1, int r, int id, Line ln ) {
      int m = (1 + r) >> 1;
      if ( nodes[ id ].id == -1 ) {
  nodes[ id ] = ln;
        return:
      bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
      if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
        atLeft ^= 1; swap( nodes[ id ], ln );
      if ( r - 1 == 1 ) return;
      if ( atLeft ) insert( 1, m, lc( id ), ln );
      else insert( m, r, rc( id ), ln );
    int query( int 1, int r, int id, int x ) {
      int ret = 0;
      if ( nodes[ id ].id != -1 )
        ret = nodes[ id ].at( x );
      int m = (1 + r) >> 1;
      if ( r - l == 1 ) return ret;
      else if ( x < m )</pre>
        return max( ret, query( 1, m, lc( id ), x ) );
      else
        return max( ret, query( m, r, rc( id ), x ) );
  public:
    void build( int n_ ) {
      n = n_; nodes.clear();
      nodes.resize( n << 2, Line() );</pre>
    void insert( Line ln ) { insert( 0, n, 0, ln ); }
    int query( int x ) { return query( 0, n, 0, x ); }
} lichao:
```

2.6 Treap

```
namespace Treap{
  #define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
  struct node{
    int size:
    uint32_t pri;
    node *lc, *rc;
    node() : size(0), pri(rand()), lc(0), rc(0) {}
    void pull() {
      size = 1
      if ( lc ) size += lc->size;
      if ( rc ) size += rc->size;
   }
 node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
    if ( L->pri > R->pri ) {
      L->rc = merge( L->rc, R ); L->pull();
      return L;
    } else {
      R->lc = merge( L, R->lc ); R->pull();
      return R;
    }
  void split_by_size( node*rt,int k,node*&L,node*&R ) {
    if ( not rt ) L = R = nullptr;
    else if( sz( rt->lc ) + 1 <= k ) {
      L = rt:
      split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
      L->pull();
    } else {
      R = rt;
```

```
split_by_size( rt->lc, k, L, R->lc );
    R->pull();
}

#undef sz
}
```

2.7 SparseTable

```
template < typename T, typename Cmp_ = less< T > >
class SparseTable {
private:
  vector< vector< T > > tbl;
  vector< int > lg;
  T cv(Ta, Tb) {
    return Cmp_()( a, b ) ? a : b;
public:
  void init( T arr[], int n ) {
     // 0-base
    lg.resize(n + 1);
    lg[0] = -1;
    for( int i=1 ; i<=n ; ++i ) lg[i] = lg[i>>1] + 1;
    tbl.resize(lg[n] + 1);
    tbl[ 0 ].resize( n );
    copy( arr, arr + n, tbl[ 0 ].begin() );
for ( int i = 1 ; i <= lg[ n ] ; ++ i ) {</pre>
       int len = 1 << ( i - 1 ), sz = 1 << i;</pre>
       tbl[ i ].resize( n - sz + 1 );
       for (int j = 0; j \le n - sz)
         tbl[i][j] = cv(tbl[i-1][j], tbl[i-1][j+len]);
  T query( int 1, int r ) {
    // 0-base [1, r)
    int wh = lg[r - 1], len = 1 << wh;
return cv( tbl[ wh ][ 1], tbl[ wh ][ r - len ] );</pre>
};
```

2.8 Linear Basis

```
struct LinearBasis {
private:
  int n, sz;
  vector< llu > B;
  inline llu two( int x ){ return ( ( llu ) 1 ) << x; }</pre>
public:
  void init( int n_ ) {
    n = n_{;} B.clear(); B.resize(n); sz = 0;
  void insert( llu x ) {
    // add x into B
    for ( int i = n-1; i >= 0 ; --i ) if( two(i) & x ){
      if (B[i]) x ^= B[i];
      else ·
        B[i] = x; sz++;
        for ( int j = i - 1 ; j >= 0 ;
          if( B[ j ] && ( two( j ) & B[ i ] ) )
        B[ i ] ^= B[ j ];
for (int j = i + 1 ;
          or (int j = i + 1 ; j < n
if ( two( i ) & B[ j ] )
                                 < n; ++ j)
            B[j] ^= B[i];
        break:
      }
    }
  inline int size() { return sz; }
  bool check( llu x ) {
    // is x in span(B) ?
    for ( int i = n-1 ; i >= 0 ; --i ) if( two(i) & x )
      if( B[ i ] ) x ^= B[ i ];
      else return false;
    return true;
  llu kth_small(llu k) {
    /** 1-base would always > 0 **/
    /** should check it **/
    /* if we choose at least one element
       but size(B)(vectors in B)==N(original elements)
       then we can't get 0 */
```

```
llu ret = 0;
for ( int i = 0 ; i < n ; ++ i ) if( B[ i ] ) {
    if( k & 1 ) ret ^= B[ i ];
    k >>= 1;
}
return ret;
}
} base;
```

3 Graph

3.1 Euler Circuit

```
bool vis[ N ]; size_t la[ K ];
void dfs( int u, vector< int >& vec ) {
  while ( la[ u ] < G[ u ].size() ) {
    if( vis[ G[ u ][ la[ u ] ].second ] ) {
        ++ la[ u ];
        continue;
    }
    int v = G[ u ][ la[ u ] ].first;
    vis[ G[ u ][ la[ u ] ].second ] = true;
    ++ la[ u ]; dfs( v, vec );
    vec.push_back( v );
  }
}</pre>
```

3.2 BCC Edge

```
class BCC{
private:
  vector< int > low, dfn;
  int cnt;
  vector< bool > bridge;
  vector< vector< PII > > G;
  void dfs( int w, int f ) {
    low[ w ] = dfn[ w ] = cnt++;
    for ( auto [ u, t ] : G[ w ] ) {
  if ( u == f ) continue;
      if ( dfn[ u ] != 0 ) {
        low[ w ] = min( low[ w ], dfn[ u ] );
        dfs( u, w );
        low[ w ] = min( low[ w ], low[ u ] );
        if ( low[ u ] > dfn[ w ] ) bridge[ t ] = true;
public:
  void init( int n, int m ) {
    G.resize(n); cnt = 0;
    fill( G.begin(), G.end(), vector< PII >() );
    bridge.clear(); bridge.resize( m );
    low.clear(); low.resize( n );
dfn.clear(); dfn.resize( n );
  void add_edge( int u, int v ) {
    // should check for multiple edge
    G[ u ].emplace_back( v, cnt );
    G[ v ].emplace_back( u, cnt ++ );
  void solve(){
    cnt = 1;
    for (int i = 0; i < n; ++i)
      if (not vis[ i ]) dfs(i, i);
  // the id will be same as insert order, 0-base
  bool is_bridge( int x ) { return bridge[ x ]; }
```

3.3 BCC Vertex

```
class BCC {
  private:
    int n, t, ecnt;
    vector<vector<pair<int, int>>> G;
    vector<int> low, tin, st, bcc;
    vector<bool> ap, ins;
    void dfs(int x, int p) {
```

```
tin[x] = low[x] = ++t;
       int ch = 0;
       for (auto u: G[x]) {
         if (u.first == p) continue;
         if (not ins[u.second]) -
           st.push_back(u.second);
           ins[u.second] = true;
         if (tin[u.first]) {
           low[x] = min(low[x], tin[u.first]);
         ++ch; dfs(u.first, x);
         low[x] = min(low[x], low[u.first]);
if (low[u.first] >= tin[x]) {
           ap[x] = true; ++ecnt;
           while (true)
             int e = st.back(); st.pop_back();
             bcc[e] = ecnt;
             if (e == u.second) break;
         }
       if (ch == 1 and p == x) ap[x] = false;
  public:
    void init(int n_) {
      n = n_, ecnt = 0; st.clear();
       G.clear(); G.resize(n);
       low.clear(); tin.clear();
ap.assign(n, false);
    void add_edge(int u, int v) {
       G[u].emplace_back(v, ecnt);
       G[v].emplace_back(u, ecnt++);
    void solve() {
       ecnt = 0; bcc.resize(t);
       ins.assign(t, false);
       for (int i = 0; i < n; ++i)
         if (low[i] == 0) dfs(i, i);
    int get_id(int x) { return bcc[x];; }
    int count() { return ecnt; }
    bool is_ap(int x) { return ap[x]; }
};
```

3.4 2-SAT (SCC)

```
class TwoSat{
  private:
    int n;
    vector<vector<int>> rG,G,sccs;
    vector<int> ord,idx;
    vector<bool> vis,result;
    void dfs(int u){
      vis[u]=true
      for(int v:G[u])
        if(!vis[v]) dfs(v);
      ord.push_back(u);
    void rdfs(int u){
      vis[u]=false;idx[u]=sccs.size()-1;
      sccs.back().push_back(u);
      for(int v:rG[u])
        if(vis[v])rdfs(v);
  public:
    void init(int n_){
      n=n_;G.clear();G.resize(n);
      rG.clear();rG.resize(n);
      sccs.clear();ord.clear();
      idx.resize(n);result.resize(n);
    void add_edge(int u,int v){
      G[u].push_back(v);rG[v].push_back(u);
    void orr(int x,int y){
      if ((x^y)==1)return;
      add_edge(x^1,y); add_edge(y^1,x);
```

```
bool solve(){
      vis.clear();vis.resize(n);
       for(int i=0;i<n;++i)</pre>
         if(not vis[i])dfs(i);
       reverse(ord.begin(),ord.end());
      for (int u:ord){
         if(!vis[u])continue;
         sccs.push_back(vector<int>());
         rdfs(u);
      for(int i=0;i<n;i+=2)</pre>
         if(idx[i]==idx[i+1])
           return false;
      vector<bool> c(sccs.size());
      for(size_t i=0;i<sccs.size();++i){</pre>
         for(size_t j=0;j<sccs[i].size();++j){</pre>
           result[sccs[i][j]]=c[i];
c[idx[sccs[i][j]^1]]=!c[i];
      }
      return true;
    bool get(int x){return result[x];}
    inline int get_id(int x){return idx[x];}
    inline int count(){return sccs.size();}
} sat2;
```

Lowbit Decomposition

```
class LowbitDecomp{
private:
  int time_, chain_, LOG_N;
  vector< vector< int > > G, fa;
 vector< int > tl, tr, chain, chain_st;
 // chain_ : number of chain
// tl, tr[ u ] : subtree interval in the seq. of u
  // chain_st[ u ] : head of the chain contains u
  // chian[ u ] : chain id of the chain u is on
  inline int lowbit( int x ) {
    return x & ( -x );
  void predfs( int u, int f ) {
    chain[ u ] = 0;
    for ( int v : G[ u ] ) {
      if ( v == f ) continue;
      predfs( v, u );
       if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
         chain[ u ] = chain[ v ];
    if ( not chain[ u ] )
      chain[ u ] = chain_ ++;
  void dfschain( int u, int f ) {
    fa[ u ][ 0 ] = f;

for ( int i = 1 ; i < LOG_N ; ++ i )

fa[ u ][ i ] = fa[ fa[ u ][ i - 1 ] ][ i - 1 ];
    tl[ u ] = time_++;
    if ( not chain_st[ chain[ u ] ] )
      chain_st[ chain[ u ] ] = u;
    for ( int v : G[ u ] )
  if ( v != f and chain[ v ] == chain[ u ] )
        dfschain( v, u );
    for ( int v : G[ u ] )
          ( v != f and chain[ v ] != chain[ u ] )
      if
         dfschain( v, u );
    tr[ u ] = time_;
  inline bool anc( int u, int v ) {
    return tl[ u ] <= tl[ v ] \
      and tr[ v ] <= tr[ u ];
public:
 inline int lca( int u, int v ) {
    if ( anc( u, v ) ) return u;
for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
       if ( not anc( fa[ u ][ i ], v ) )
        u = fa[ u ][ i ];
    return fa[ u ][ 0 ];
  void init( int n ) {
    n ++;
```

```
for (LOG_N = 0 ; (1 << LOG_N) < n ; ++ LOG_N);
     fa.clear();
     fa.resize( n, vector< int >( LOG_N ) );
    G.clear(); G.resize( n );
tl.clear(); tl.resize( n );
tr.clear(); tr.resize( n );
     chain.clear(); chain.resize( n );
     chain_st.clear(); chain_st.resize( n );
  void add_edge( int u , int v ) {
     G[ u ].push_back( v );
     G[ v ].push_back( u );
  void decompose(){
     chain_ = 1;
     predfs( 1, 1 );
     time_{-} = 0;
     dfschain(1,1);
  PII get_inter( int u ) { return {tl[ u ], tr[ u ]}; }
  vector< PII > get_path( int u , int v ){
  vector< PII > res;
     int g = lca( u, v );
     while ( chain[ u ] != chain[ g ] ) {
       int s = chain_st[ chain[ u ] ]
       res.emplace_back( tl[ s ], tl[ u ] + 1 );
       u = fa[ s ][ 0 ];
    res.emplace_back( tl[ g ], tl[ u ] + 1 );
while ( chain[ v ] != chain[ g ] ) {
  int s = chain_st[ chain[ v ] ];
       res.emplace_back( tl[ s ], tl[ v ] + 1 );
       v = fa[ s ][ 0 ];
     res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
     return res
     /* res : list of intervals from u to v
      \star ( note only nodes work, not edge )
      * vector< PII >& path = tree.get_path( u , v )
      * for( auto [ 1, r ] : path ) {
          0-base [ 1, r )
      * }
      */
} tree;
```

3.6 MaxClique

```
// contain a self loop u to u, than u won't in clique
template < size_t MAXN >
class MaxClique{
private:
  using bits = bitset< MAXN >;
  bits popped, G[ MAXN ], ans;
  size_t deg[ MAXN ], deo[ MAXN ], n;
  void sort_by_degree() {
    popped.reset();
     for ( size_t i = 0 ; i < n ; ++ i )
    deg[ i ] = G[ i ].count();
for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
         size_t mi = MAXN, id = 0;
             ( size_t j = 0 ; j < n ; ++ j ) if ( not popped[ j ] and deg[ j ] < mi )
         for ( size_t j = 0 ;
                  mi = deg[id = j];
         popped[ deo[ i ] = id ] = 1;
for( size_t u = G[ i ]._Find_first()
           u < n ; u = G[ i ]._Find_next( u ) )
              -- deg[ u ];
    }
  void BK( bits R, bits P, bits X ) {
    if (R.count()+P.count() <= ans.count()) return;</pre>
    if ( not P.count() and not X.count() ) {
       if ( R.count() > ans.count() ) ans = R;
       return;
     /* greedily chosse max degree as pivot
    bits cur = P | X; size_t pivot = 0, sz = 0;
    for ( size_t u = cur._Find_first() ;
```

```
u < n ; u = cur._Find_next( u ) )
if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
    cur = P & ( ~G[ pivot ] );
     */ // or simply choose first
    bits cur = P & (~G[ ( P | X )._Find_first() ]);
    for ( size_t u = cur._Find_first() ;
       u < n ; u = cur._Find_next( u ) ) {
if ( R[ u ] ) continue;</pre>
       R[u] = 1;
       BK( R, P & G[ u ], X & G[ u ] );
       R[u] = P[u] = 0, X[u] = 1;
    }
public:
  void init( size_t n_ ) {
    n = n_{-};
    for ( size_t i = 0 ; i < n ; ++ i )
       G[ i ].reset();
    ans.reset();
  void add_edges( int u, bits S ) { G[ u ] = S; }
  void add_edge( int u, int v ) {
    G[u][v] = G[v][u] = 1;
  int solve() {
    sort_by_degree(); // or simply iota( deo... )
for ( size_t i = 0 ; i < n ; ++ i )</pre>
       deg[ i ] = G[ i ].count();
    bits pob, nob = 0; pob.set();
    for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
    for ( size_t i = 0 ; i < n ; ++ i ) {
       size_t v = deo[ i ];
       bits tmp; tmp[ v ] = 1;
BK( tmp, pob & G[ v ], nob & G[ v ] );
pob[ v ] = 0, nob[ v ] = 1;
     return static_cast< int >( ans.count() );
};
```

3.7 Virtural Tree

```
inline bool cmp(const int &i, const int &j) {
  return dfn[i] < dfn[j];</pre>
void build(int vectrices[], int k) {
 static int stk[MAX_N];
  sort(vectrices, vectrices + k, cmp);
  stk[sz++] = 0;
  for (int i = 0; i < k; ++i) {
  int u = vectrices[i], lca = LCA(u, stk[sz - 1]);</pre>
    if (lca == stk[sz - 1]) stk[sz++] = u;
      while (sz \ge 2 \&\& dep[stk[sz - 2]] \ge dep[lca]) {
        addEdge(stk[sz - 2], stk[sz - 1]);
        sz--;
      if (stk[sz - 1] != lca) {
        addEdge(lca, stk[--sz]);
        stk[sz++] = lca, vectrices[cnt++] = lca;
      stk[sz++] = u;
    }
  for (int i = 0; i < sz - 1; ++i)
    addEdge(stk[i], stk[i + 1]);
```

3.8 Tree Hashing

```
uint64_t hsah( int u, int f ) {
    uint64_t r = 127;
    for ( int v : G[ u ] ) {
        if ( v == f ) continue;
        uint64_t hh = hsah( v, u );
        r = r + ( hh * hh ) % mod;
    }
    return r;
}
```

3.9 Minimum Mean Cycle

```
/* minimum mean cycle O(VE) */
struct MMC{
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
#define V 1021
#define inf 1e9
   struct Edge { int v,u; double c; };
   int n, m, prv[V][V], prve[V][V], vst[V];
   Edge e[E];
   vector<int> edgeID, cycle, rho;
   double d[V][V];
  void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
   void add_edge( int vi , int ui , double ci )
  { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
     for(int i=0; i<n; i++) d[0][i]=0;</pre>
     for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);</pre>
       for(int j=0; j<m; j++)</pre>
         int v = e[j].v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
            d[i+1][u] = d[i][v]+e[j].c;
            prv[i+1][u] = v;
            prve[i+1][u] = j;
         }
       }
    }
   double solve(){
     // returns inf if no cycle, mmc otherwise
     double mmc=inf;
     int st = -1
     bellman_ford();
     for(int i=0; i<n; i++) {</pre>
       double avg=-inf;
       for(int k=0; k<n; k++) {
  if(d[n][i]<inf-eps)</pre>
            avg=max(avg,(d[n][i]-d[k][i])/(n-k));
          else avg=max(avg,inf);
       if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
     FZ(vst);edgeID.clear();cycle.clear();rho.clear();
     for (int i=n; !vst[st]; st=prv[i--][st]) {
       vst[st]++
       edgeID.PB(prve[i][st]);
       rho.PB(st);
     while (vst[st] != 2) {
       int v = rho.back(); rho.pop_back();
       cycle.PB(v);
       vst[v]++;
     reverse(ALL(edgeID));
     edgeID.resize(SZ(cycle));
     return mmc;
} mmc;
3.10 Mo's Algorithm on Tree
```

```
int q; vector< int > G[N];
struct Que{
  int u, v, id;
} que[ N ];
int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
void dfs( int u, int f ) {
  dfn[ u ] = dfn_++; int saved_rbp = stk_;
  for ( int v : G[ u ] ) {
    if ( v == f ) continue;
    dfs( v, u );
    if ( stk_ - saved_rbp < SQRT_N ) continue;
    for ( ++ block_ ; stk_ != saved_rbp ; )
        block_id[ stk[ -- stk_ ] ] = block_;
}
stk[ stk_ ++ ] = u;
}
bool inPath[ N ];</pre>
```

```
void Diff( int u ) {
  if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
  else { /*add this edge*/ }
void traverse( int& origin_u, int u ) {
  for ( int g = lca( origin_u, u ) ;
    origin_u != g ; origin_u = parent_of[ origin_u ] )
      Diff( origin_u );
  for (int v = u; v != origin_u; v = parent_of[v])
    Diff( v );
  origin_u = u;
void solve() {
 dfs( 1, 1 );
while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
  sort( que, que + q, [](const Que& x, const Que& y) {
    return tie( block_id[ x.u ], dfn[ x.v ] )
            < tie( block_id[ y.u ], dfn[ y.v ] );
  int U = 1, V = 1;
  for ( int i = 0 ; i < q ; ++ i ) {
    pass( U, que[ i ].u );
    pass( V, que[ i ].v );
    // we could get our answer of que[ i ].id
}
/*
Method 2:
dfs u:
 push u
  iterate subtree
  push u
Let P = LCA(u, v), and St(u) <= St(v)
if (P == u) query[St(u), St(v)]
else query[Ed(u), St(v)], query[St(P), St(P)]
```

3.11 Minimum Steiner Tree

```
// Minimum Steiner Tree
// 0(V 3^T + V^2 2^T)
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
  int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
  void init( int _n ){
    n = _n;
for( int i = 0 ; i < n ; i ++ ){
       for( int j = 0 ; j < n ; j ++ )</pre>
       dst[ i ][ j ] = INF;
dst[ i ][ i ] = 0;
    }
  void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
    dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
  void shortest_path(){
    for( int k = 0 ; k < n ; k ++ )
       for( int i = 0 ; i < n ; i ++ )</pre>
         for( int j = 0 ; j < n ; j ++ )
            dst[ i ][ j ] = min( dst[ i ][ j ],
                   dst[i][k]+dst[k][j]);
  int solve( const vector<int>& ter ){
    int t = (int)ter.size();
for( int i = 0 ; i < ( 1 << t ) ; i ++ )</pre>
       for( int j = 0 ; j < n ; j ++ )
  dp[ i ][ j ] = INF;</pre>
    for( int i = 0 ; i < n ; i ++ )</pre>
       d\hat{p}[0][i] = 0;
    for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
       if( msk == ( msk & (-msk) ) ){
         int who = __lg( msk );
for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];</pre>
         continue:
       for( int i = 0 ; i < n ; i ++ )</pre>
         for( int submsk = ( msk - 1 ) & msk ; submsk ;
```

3.12 Directed Minimum Spanning Tree

```
template <typename T> struct DMST {
  T g[maxn][maxn], fw[maxn];
  int n, fr[maxn];
  bool vis[maxn], inc[maxn];
  void clear() {
    for(int i = 0; i < maxn; ++i) {</pre>
      for(int j = 0; j < maxn; ++j) g[i][j] = inf;
      vis[i] = inc[i] = false;
  }
  void addEdge(int u,int v,T w){g[u][v]=min(g[u][v],w)
      ;}
  T operator()(int root, int _n) {
    n = n; T ans = 0;
    if (dfs(root) != n) return -1;
    while (true) {
      for(int i = 1;i <= n;++i) fw[i] = inf, fr[i] = i;</pre>
      for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
        for (int j = 1; j <= n; ++j) {
          if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
             fw[i] = g[j][i]; fr[i] = j;
          }
        }
      int x = -1;
      for(int i = 1;i <= n;++i)if(i != root && !inc[i])</pre>
        int j = i, c = 0;
        while(j!=root && fr[j]!=i && c<=n) ++c, j=fr[j
        if (j == root || c > n) continue;
        else { x = i; break; }
      if (!~x) {
        for (int i = 1; i <= n; ++i)</pre>
          if (i != root && !inc[i]) ans += fw[i];
        return ans:
      int y = x;
for (int i = 1; i <= n; ++i) vis[i] = false;
        ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true
      } while (y != x);
      inc[x] = false;
      for (int k = 1; k <= n; ++k) if (vis[k])</pre>
        for (int j = 1; j <= n; ++j) if (!vis[j]) {
          if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
          if (g[j][k] < inf && g[j][k]-fw[k] < g[j][x])
             g[j][x] = g[j][k] - fw[k];
        }
      }
    }
    return ans;
  int dfs(int now) {
    int r = 1; vis[now] = true;
    for (int i = 1; i <= n; ++i)
  if (g[now][i] < inf && !vis[i]) r += dfs(i);</pre>
```

return r:

```
Dominator Tree
3.13
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
  // vertices are numbered from 0 to n - 1
  fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
  fill(fa, fa + n, -1); fill(val, val + n, -1);
  fill(sdom, sdom + n, -1); fill(rp, rp + n, -1); fill(dom, dom + n, -1); tk = \theta;
  for (int i = 0; i < n; ++i) {
    g[i].clear(); r[i].clear(); rdom[i].clear();
  }
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
  rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
  for (int u : g[x]) {
    if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
    r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
 if (fa[x] == x) return c ? -1 : x;
  int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
  if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
  fa[x] = p;
  return c ? p : val[x];
vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
// p[i] = -2 if i is unreachable from s
  dfs(s);
  for (int i = tk - 1; i >= 0; --i) {
    for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)])
    if (i) rdom[sdom[i]].push_back(i);
    for (int &u : rdom[i]) {
      int p = find(u);
      if (sdom[p] == i) dom[u] = i;
      else dom[u] = p;
    if (i) merge(i, rp[i]);
  vector<int> p(n, -2); p[s] = -1;
  for (int i = 1; i < tk; ++i)</pre>
    if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
  for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
  return p;
4
     Matching & Flow
```

4.1 Kuhn Munkres

```
class KM {
private:
    static constexpr lld INF = 1LL << 60;
    vector<lld> hl,hr,slk;
    vector<int> fl,fr,pre,qu;
    vector<vector<lld> w;
    vector<bool> vl,vr;
    int n, ql, qr;
    bool check(int x) {
        if (vl[x] = true, fl[x]! = -1)
            return vr[qu[qr++] = fl[x]] = true;
        while (x! = -1) swap(x, fr[fl[x] = pre[x]]);
        return false;
    }
    void bfs(int s) {
        fill(slk.begin(), slk.end(), INF);
    }
}
```

```
fill(v1.begin(), v1.end(), false);
    fill(vr.begin(), vr.end(), false);
    ql = qr = 0;
    qu[qr++] = s;
    vr[s] = true;
    while (true) {
      11d d;
      while (ql < qr) {</pre>
        for (int x = 0, y = qu[ql++]; x < n; ++x)
          if(!v1[x]\&s1k[x]>=(d=h1[x]+hr[y]-w[x][y])){
             if (pre[x] = y, d) slk[x] = d;
             else if (!check(x)) return;
          }
        }
      }
      d = INF;
      for (int x = 0; x < n; ++x)
  if (!v1[x] && d > s1k[x]) d = s1k[x];
       for (int x = 0; x < n; ++x) {
        if (v1[x]) h1[x] += d;
         else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
         if (!v1[x] && !slk[x] && !check(x)) return;
public:
  void init( int n_ ) {
    n = n_; qu.resize(n);
    fl.clear(); fl.resize(n, -1);
    fr.clear(); fr.resize(n, -1);
    hr.clear(); hr.resize(n); hl.resize(n);
    w.clear(); w.resize(n, vector<lld>(n));
    slk.resize(n); pre.resize(n);
    vl.resize(n); vr.resize(n);
  void set_edge( int u, int v, lld x ) {w[u][v] = x;}
  11d solve() {
    for (int i = 0; i < n; ++i)
      hl[i] = *max_element(w[i].begin(), w[i].end());
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    11d res = 0;
    for (int i = 0; i < n; ++i) res += w[i][f1[i]];</pre>
    return res;
} km;
```

4.2 Bipartite Matching

```
class BipartiteMatching{
private:
  vector<int> X[N], Y[N];
  int fX[N], fY[N], n;
  bitset<N> walked;
  bool dfs(int x){
    for(auto i:X[x]){
      if(walked[i])continue;
      walked[i]=1;
      if(fY[i]==-1||dfs(fY[i])){
        fY[i]=x;fX[x]=i;
        return 1;
      }
    }
    return 0;
public:
  void init(int _n){
    n=_n; walked.reset();
    for(int i=0;i<n;i++){</pre>
      X[i].clear();Y[i].clear();
      fX[i]=fY[i]=-1;
  void add_edge(int x, int y){
    X[x].push_back(y); Y[y].push_back(y);
  int solve(){
    int cnt = 0;
    for(int i=0;i<n;i++){</pre>
      walked.reset();
```

```
if(dfs(i)) cnt++;
    // return how many pair matched
    return cnt:
};
```

4.3 General Graph Matching

```
const int N = 514, E = (2e5) * 2;
struct Graph{
  int to[E],bro[E],head[N],e;
  int lnk[N], vis[N], stp, n;
  void init( int _n ){
  stp = 0; e = 1; n = _n;
    for( int i = 0 ; i <= n ;</pre>
                                i ++ )
      head[i] = lnk[i] = vis[i] = 0;
  void add_edge(int u,int v){
    to[e]=v,bro[e]=head[u],head[u]=e++;
    to[e]=u,bro[e]=head[v],head[v]=e++;
  bool dfs(int x){
    vis[x]=stp;
    for(int i=head[x];i;i=bro[i]){
      int v=to[i];
      if(!lnk[v]){
        lnk[x]=v, lnk[v]=x;
         return true
      }else if(vis[lnk[v]]<stp){</pre>
         int w=lnk[v]
         lnk[x]=v, lnk[v]=x, lnk[w]=0;
         if(dfs(w)) return true
        lnk[w]=v, lnk[v]=w, lnk[x]=0;
      }
    }
    return false;
  int solve(){
    int ans = 0;
    for(int i=1;i<=n;i++)</pre>
      if(not lnk[i]){
        stp++; ans += dfs(i);
    return ans;
} graph;
```

4.4 Minimum Weight Matching (Clique version)

```
struct Graph {
  // 0-base (Perfect Match)
  int n, edge[MXN][MXN];
  int match[MXN], dis[MXN], onstk[MXN];
  vector<int> stk;
  void init(int _n) {
    n = _n;
    for (int i=0; i<n; i++)</pre>
      for (int j=0; j<n; j++)</pre>
        edge[i][j] = 0;
  void set_edge(int u, int v, int w) {
    edge[u][v] = edge[v][u] = w;
  bool SPFA(int u){
    if (onstk[u]) return true;
    stk.PB(u);
    onstk[u] = 1;
    for (int v=0; v<n; v++){</pre>
      if (u != v && match[u] != v && !onstk[v]){
        int m = match[v];
        if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
          dis[m] = dis[u] - edge[v][m] + edge[u][v];
          onstk[v] = 1;
          stk.PB(v);
          if (SPFA(m)) return true;
          stk.pop_back();
          onstk[v] = 0;
      }
```

```
onstk[u] = 0;
    stk.pop_back();
    return false;
  int solve() {
    // find a match
    for (int i=0; i<n; i+=2){</pre>
      match[i] = i+1;
      match[i+1] = i;
    while (true){
  int found = 0;
       for (int i=0; i<n; i++)</pre>
         dis[i] = onstk[i] = 0;
       for (int i=0; i<n; i++){</pre>
         stk.clear(
         if (!onstk[i] && SPFA(i)){
           found = 1;
           while (SZ(stk)>=2){
             int u = stk.back(); stk.pop_back();
             int v = stk.back(); stk.pop_back();
             match[u] = v;
             match[v] = u;
           }
        }
      if (!found) break;
    int ret = 0;
    for (int i=0; i<n; i++)</pre>
      ret += edge[i][match[i]];
    return ret>>1;
} graph;
```

4.5 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}$.
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect
 - To maximize, connect t o s with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f
 eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the
 - maximum flow from s to t is the answer.

 To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f'
 eq \sum_{v \in V, in(v)>0}^{\cdot} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- $\bullet\,$ Construct minimum vertex cover from maximum matching M on bipartite graph(X,Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X. 3. $x \in X$ is chosen iff x is unvisited.

 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect $x \to y$ with (cost,cap) = (c,1) if c>0, otherwise connect $y\to x$ with (cost, cap)=(-c,1)
 - 3. For each edge with c < 0, sum these cost as K, then increase d(y)by 1, decrease $d(\boldsymbol{x})$ by 1
 - 4. For each vertex v with d(v)>0, connect S o v with (cost, cap)=(0, d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =(0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \to v$, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity
 - 5. For $v \in G$, connect it with sink $v \to t$ with capacity K + 2T - $(\sum_{e \in E(v)} w(e)) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K |V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u, v).

- 2. Connect $v \, o \, v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
- 3. Find the minimum weight perfect matching on G^{\prime} .
- Project selection problem
 - 1. If $p_{v}>0$, create edge $\left(s,v\right)$ with capacity p_{v} ; otherwise, create edge
 - (v,t) with capacity $-p_v$. 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

```
\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y'})
```

can be minimized by the mincut of the following graph:

- 1. Create edge (x, t) with capacity c_x and create edge (s, y) with ca-
- 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

Dinic 4.6

```
class Dinic{
private:
  using CapT = int64_t;
  struct Edge{
    int to, rev;
    CapT cap;
  int n, st, ed;
  vector<vector<Edge>> G;
  vector<int> lv, idx;
  bool BFS(){
    fill(lv.begin(), lv.end(), -1);
    queue<int> bfs;
    bfs.push(st);
    lv[st] = 0;
    while(!bfs.empty()){
      int u = bfs.front(); bfs.pop();
      for(auto e: G[u]){
        if(e.cap <= 0 or lv[e.to]!=-1) continue;
lv[e.to] = lv[u] + 1;</pre>
        bfs.push(e.to);
      }
    return (lv[ed]!=-1);
  CapT DFS(int u, CapT f){
    if(u == ed) return f;
    CapT ret = 0;
    for(int& i = idx[u]; i < (int)G[u].size(); ++i){</pre>
      auto& e = G[u][i]
      if(e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
      CapT nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
      G[e.to][e.rev].cap += nf;
      if(f == 0) return ret;
    if(ret == 0) lv[u] = -1;
    return ret;
public:
  void init(int n_, int st_, int ed_){
    n = n_, st = st_, ed = ed_;
    G.resize(n); lv.resize(n);
    fill(G.begin(), G.end(), vector<Edge>());
  void add_edge(int u, int v, CapT c){
    G[u].push_back({v, (int)G[v].size(), c});
    G[v].push_back({u, ((int)G[u].size())-1, 0});
  CapT max_flow(){
    CapT ret = 0;
    while(BFS()){
      idx.assign(n, 0);
      CapT f = DFS(st, numeric_limits<CapT>::max());
      ret += f;
      if(f == 0) break;
    }
    return ret;
} flow;
```

4.7 Minimum Cost Maximum Flow

```
class MiniCostMaxiFlow{
  using CapT = int;
  using WeiT = int64_t;
  using PCW = pair<CapT,WeiT>;
  static constexpr CapT INF_CAP = 1 << 30;</pre>
  static constexpr WeiT INF_WEI = 1LL<<60;</pre>
private:
  struct Edge{
    int to, back;
    WeiT wei;
    CapT cap;
    Edge() {}
    Edge(int a,int b,WeiT c,CapT d):
      to(a),back(b),wei(c),cap(d)
    {}
  int ori, edd;
  vector<vector<Edge>> G;
  vector<int> fa, wh;
  vector<bool> inq;
  vector<WeiT> dis;
  PCW SPFA(){
    fill(inq.begin(),inq.end(),false);
    fill(dis.begin(), dis.end(), INF_WEI);
    queue<int> qq; qq.push(ori);
    dis[ori]=0;
    while(!qq.empty()){
      int_u=qq.front();qq.pop();
      inq[u] = 0;
      for(int i=0;i<SZ(G[u]);++i){</pre>
        Edge e=G[u][i];
        int v=e.to;
        WeiT d=e.wei;
        if(e.cap<=0||dis[v]<=dis[u]+d)</pre>
          continue
        dis[v]=dis[u]+d;
        fa[v]=u,wh[v]=i;
        if(inq[v]) continue;
        qq.push(v);
        inq[v]=1;
    if(dis[edd]==INF_WEI)
      return {-1,-1};
    CapT mw=INF_CAP;
    for(int i=edd;i!=ori;i=fa[i])
      mw=min(mw,G[fa[i]][wh[i]].cap);
    for (int i=edd;i!=ori;i=fa[i]){
      auto &eg=G[fa[i]][wh[i]];
      eg.cap-=mw;
      G[eg.to][eg.back].cap+=mw;
    return {mw,dis[edd]};
public:
  void init(int a,int b,int n){
    ori=a,edd=b;
    G.clear();G.resize(n);
    fa.resize(n);wh.resize(n);
    inq.resize(n); dis.resize(n);
  void add_edge(int st,int ed,WeiT w,CapT c){
    G[st].emplace_back(ed,SZ(G[ed]),w,c);
    G[ed].emplace_back(st,SZ(G[st])-1,-w,0);
  PCW solve(){
    /* might modify to
    cc += ret.first * ret.second
    or
    ww += ret.first * ret.second
    CapT cc=0; WeiT ww=0;
    while(true){
      PCW ret=SPFA();
      if(ret.first==-1) break;
      cc+=ret.first;
      ww+=ret.second;
    return {cc,ww};
```

```
National Taiwan University - kiseki
|} mcmf;
 4.8
        Global Min-Cut
 const int maxn = 500 + 5;
 int w[maxn][maxn], g[maxn];
 bool v[maxn], del[maxn];
 void add_edge(int x, int y, int c) {
   w[x][y] += c; w[y][x] += c;
 pair<int, int> phase(int n) {
   memset(v, false, sizeof(v));
   memset(g, 0, sizeof(g));
int s = -1, t = -1;
   while (true) {
      int c = -1;
      for (int i = 0; i < n; ++i) {
         if (del[i] || v[i]) continue;
         if (c == -1 \mid | g[i] > g[c]) c = i;
      if (c == -1) break;
      v[s = t, t = c] = true;
      for (int i = 0; i < n; ++i) {
        if (del[i] || v[i]) continue;
         g[i] += w[c][i];
      }
   return make_pair(s, t);
 int mincut(int n) {
   int cut = 1e9;
   memset(del, false, sizeof(del));
   for (int i = 0; i < n - 1; ++i) {
      int s, t; tie(s, t) = phase(n);
del[t] = true; cut = min(cut, g[t]);
for (int j = 0; j < n; ++j) {</pre>
         w[s][j] += w[t][j]; w[j][s] += w[j][t];
   return cut;
 5
       Math
       Prime Table
 1002939109, 1020288887, 1028798297, 1038684299,
 1041211027, 1051762951, 1058585963, 1063020809,
 1147930723, 1172520109, 1183835981, 1187659051,\\
 1241251303, 1247184097, 1255940849, 1272759031,\\
 \begin{array}{c} 1287027493, 1288511629, 1294632499, 1312650799, \\ 1868732623, 1884198443, 1884616807, 1885059541, \\ 1909942399, 1914471137, 1923951707, 1925453197, \end{array}
 1979612177, 1980446837, 1989761941, 2007826547,
 2008033571, 2011186739, 2039465081, 2039728567
 2093735719, 2116097521, 2123852629, 2140170259,\\
 3148478261, 3153064147, 3176351071, 3187523093,
 3196772239, 3201312913, 3203063977, 3204840059, 3210224309, 3213032591, 3217689851, 3218469083,
 3219857533, 3231880427, 3235951699, 3273767923,
 3276188869, 3277183181, 3282463507, 3285553889,
 3319309027, 3327005333, 3327574903, 3341387953,
 3373293941, 3380077549, 3380892997, 3381118801
 5.2 \lfloor \frac{n}{i} \rfloor Enumeration
 T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_{i+1} \rfloor} \rfloor} \rfloor
 5.3 ax+by=gcd
 // ax+ny = 1, ax+ny == ax == 1 \pmod{n}
 void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
   if (y == 0) g=x, a=1, b=0;
   else exgcd(y,x%y,g,b,a),b=(x/y)*a;
         Pollard Rho
 // does not work when n is prime
 // return any non-trivial factor
 llu pollard_rho(llu n){
   static auto f=[](llu x,llu k,llu m){
```

return add(k,mul(x,x,m),m);

```
mt19937 rnd(120821011);
  while(true){
    llu y=2, yy=y, x=rnd()%n, t=1;
    for(llu sz=2;t==1;sz<<=1) {</pre>
      for(llu i=0;i<sz;++i){</pre>
        if(t!=1)break;
        yy=f(yy,x,n);
        t=gcd(yy>y?yy-y:y-yy,n);
      y=yy;
    if(t!=1&&t!=n) return t;
}
5.5
      Pi Count (Linear Sieve)
static constexpr int N = 1000000 + 5;
lld pi[N];
vector<int> primes;
bool sieved[N];
11d cube_root(11d x){
  1ld s=cbrt(x-static_cast<long double>(0.1));
  while(s*s*s \ll x) ++s;
  return s-1;
1ld square_root(1ld x){
  1ld s=sqrt(x-static_cast<long double>(0.1));
  while(s*s <= x) ++s;</pre>
  return s-1;
void init(){
  primes.reserve(N);
  primes.push_back(1);
  for(int i=2;i<N;i++) {
    if(!sieved[i]) primes.push_back(i);
    pi[i] = !sieved[i] + pi[i-1];
    for(int p: primes) if(p > 1) {
      if(p * i >= N) break;
      sieved[p * i] = true;
      if(p % i == 0) break;
  }
1ld phi(lld m, lld n) {
  static constexpr int MM = 80000, NN = 500;
  static lld val[MM][NN];
  if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
  if(n == 0) return m;
  if(primes[n] >= m) return 1;
  11d ret = phi(m,n-1)-phi(m/primes[n],n-1);
  if(m < MM\&n < NN) val[m][n] = ret+1;
  return ret;
11d pi_count(11d);
11d P2(11d m, 11d n) {
  11d sm = square_root(m), ret = 0;
  for(lld i = n+1;primes[i]<=sm;i++)</pre>
    ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
  return ret;
11d pi_count(11d m) {
  if(m < N) return pi[m];</pre>
  11d n = pi_count(cube_root(m));
  return phi(m, n) + n - 1 - P2(m, n);
5.6 Range Sieve
const int MAX_SQRT_B = 50000;
const int MAX_L = 200000 + 5;
bool is_prime_small[MAX_SQRT_B];
bool is_prime[MAX_L];
void sieve(lld 1, lld r){
  // [1, r)
  for(lld i=2;i*i<r;i++) is_prime_small[i] = true;</pre>
```

for(lld i=l;i<r;i++) is_prime[i-l] = true;</pre>

```
if(l==1) is_prime[0] = false;
for(lld i=2;i*i<r;i++){
    if(!is_prime_small[i]) continue;
    for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;
    for(lld j=std::max(2LL, (1+i-1)/i)*i;j<r;j+=i)
        is_prime[j-1]=false;
}
</pre>
```

5.7 Miller Rabin

```
bool isprime(llu x){
 static llu magic[]={2,325,9375,28178,\
                     450775,9780504,1795265022};
  static auto witn=[](llu a,llu u,llu n,int t){
   a = mpow(a,u,n);
    if (!a)return 0;
    while(t--){
      11u a2=mul(a,a,n);
      if(a2==1 && a!=1 && a!=n-1)
        return 1;
      a = a2:
    }
    return a!=1;
  if(x<2)return 0;</pre>
  if(!(x&1))return x==2;
 llu x1=x-1; int t=0;
  while(!(x1&1))x1>>=1,t++;
 for(llu m:magic)if(witn(m,x1,x,t))return 0;
 return 1:
```

5.8 Inverse Element

```
// x's inverse mod k
long long GetInv(long long x, long long k){
  // k is prime: euler_(k)=k-1
  return qPow(x, euler_phi(k)-1);
}

// if you need [1, x] (most use: [1, k-1]
void solve(int x, long long k){
  inv[1] = 1;
  for(int i=2;i<x;i++)
    inv[i] = ((long long)(k - k/i) * inv[k % i]) % k;
}</pre>
```

5.9 Euler Phi Function

```
extended euler:
   a^b mod p
   if gcd(a, p)==1: a^(b%phi(p))
   elif b < phi(p): a^b mod p
   else a^(b%phi(p) + phi(p))
lld euler_phi(int x){
 11d r=1;
  for(int i=2;i*i<=x;++i){</pre>
    if(x\%i==0){
      x/=i; r*=(i-1);
      while(x%i==0) {
        x/=i; r*=i;
   }
  if(x>1) r*=x-1;
  return r;
vector<int> primes;
bool notprime[N];
11d phi[N];
void euler_sieve(int n){
  for(int i=2;i<n;i++){</pre>
    if(!notprime[i]){
      primes.push_back(i); phi[i] = i-1;
    for(auto j: primes){
      if(i*j >= n) break;
      notprime[i*j] = true;
```

```
phi[i*j] = phi[i] * phi[j];
if(i % j == 0) {
    phi[i*j] = phi[i] * j;
    break;
}
}
}
}
```

5.10 Gauss Elimination

```
void gauss(vector<vector<double>> &d) {
   int n = d.size(), m = d[0].size();
   for (int i = 0; i < m; ++i) {
      int p = -1;
      for (int j = i; j < n; ++j) {
        if (fabs(d[j][i]) < eps) continue;
        if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
      }
   if (p == -1) continue;
   for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
   for (int j = 0; j < n; ++j) {
      if (i == j) continue;
      double z = d[j][i] / d[i][i];
      for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
   }
}</pre>
```

5.11 Fast Fourier Transform

```
polynomial multiply:
   DFT(a, len); DFT(b, len);
   for(int i=0;i<len;i++) c[i] = a[i]*b[i];
   iDFT(c, len);
   (len must be 2^k and = 2^k(max(a, b)))
   Hand written Cplx would be 2x faster
Cplx omega[2][N];
void init_omega(int n) -
  static constexpr llf PI=acos(-1);
  const llf arg=(PI+PI)/n;
  for(int i=0;i<n;++i)</pre>
    omega[0][i]={cos(arg*i), sin(arg*i)};
  for(int i=0;i<n;++i)</pre>
    omega[1][i]=conj(omega[0][i]);
void tran(Cplx arr[],int n,Cplx omg[]) {
  for(int i=0, j=0;i<n;++i){</pre>
    if(i>j)swap(arr[i],arr[j]);
    for(int l=n>>1;(j^=1)<1;l>>=1);
  for (int l=2;l<=n;l<<=1){
    int m=1>>1;
    for(auto p=arr;p!=arr+n;p+=1){
      for(int i=0;i<m;++i)</pre>
        Cplx t=omg[n/1*i]*p[m+i];
        p[m+i]=p[i]-t; p[i]+=t;
      }
    }
  }
void DFT(Cplx arr[],int n){tran(arr,n,omega[0]);}
void iDFT(Cplx arr[],int n){
  tran(arr,n,omega[1]);
  for(int i=0;i<n;++i) arr[i]/=n;</pre>
```

5.12 High Speed Linear Recurrence

```
#define mod 998244353
const int N=1000010;
int n,k,m,f[N],h[N],a[N],b[N],ib[N];
int pw(int x,int y){
  int re=1;
  if(y<0)y+=mod-1;
  while(y){
   if(y&1)re=(l1)re*x%mod;
  y>>=1;x=(l1)x*x%mod;
```

```
reverse(a,a+k+1);
  return re;
                                                                 poly::inv(a,ib,len);
                                                                 poly::cls(ib,k+1,len);
void inc(int&x,int y){x+=y;if(x>=mod)x-=mod;}
                                                                 poly::ntt(b,len,1);
                                                                 poly::ntt(ib,len,1);
namespace poly{
 const int G=3;
                                                                 poly::pow(a,n);
  int rev[N],L;
                                                                 int ans=0;
  void ntt(int*A,int len,int f){
                                                                 for(int i=0;i<k;++i)inc(ans,(ll)a[i]*h[i]%mod);</pre>
                                                                 printf("%d\n",ans);
    for(L=0;(1<<L)<len;++L);</pre>
    for(int i=0;i<len;++i){</pre>
                                                                 return 0;
      rev[i]=(rev[i>>1]>>1)|((i&1)<<(L-1));
      if(i<rev[i])swap(A[i],A[rev[i]]);</pre>
                                                               5.13
                                                                      Chinese Remainder
    for(int i=1;i<len;i<<=1){</pre>
      int wn=pw(G, f*(mod-1)/(i<<1));</pre>
                                                               lld crt(lld ans[], lld pri[], int n){
      for(int j=0;j<len;j+=i<<1){</pre>
                                                                 11d M = 1, ret = 0;
        int w=1;
                                                                 for(int i=0;i<n;i++) M *= pri[i];</pre>
        for(int k=0;k<i;++k,w=(11)w*wn%mod){</pre>
                                                                 for(int i=0;i<n;i++){</pre>
          int x=A[j+k],y=(11)w*A[j+k+i]%mod;
                                                                   1ld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
          A[j+k]=(x+y)\mbox{mod}, A[j+k+i]=(x-y+mod)\mbox{mod};
                                                                   ret += (ans[i]*(M/pri[i])%M * iv)%M;
                                                                   ret %= M;
      }
                                                                 return ret;
    if(!~f){
                                                               }
      int iv=pw(len,mod-2);
                                                               /*
      for(int i=0;i<len;++i)A[i]=(11)A[i]*iv%mod;</pre>
                                                               Another:
                                                               x = a1 \% m1
  }
                                                               x = a2 \% m2
  void cls(int*A,int l,int r)
                                                               g = gcd(m1, m2)
    for(int i=1;i<r;++i)A[i]=0;}</pre>
                                                               assert((a1-a2)%g==0)
  void cpy(int*A,int*B,int 1){
                                                               [p, q] = exgcd(m2/g, m1/g)
    for(int i=0;i<1;++i)A[i]=B[i];}</pre>
                                                               return a2+m2*(p*(a1-a2)/g)
  void inv(int*A,int*B,int 1){
                                                               0 <= x < lcm(m1, m2)
    if(l==1){B[0]=pw(A[0],mod-2);return;}
                                                               */
    static int t[N];
    int len=l<<1;
    inv(A,B,l>>1);
                                                               5.14
                                                                      Berlekamp Massey
    cpy(t, A, 1); cls(t, 1, len);
    ntt(t,len,1);ntt(B,len,1);
                                                               // x: 1-base, p[]: 0-base
    for(int i=0;i<len;++i)</pre>
                                                               template<size_t N>
      B[i]=(11)B[i]*(2-(11)t[i]*B[i]*mod+mod)*mod;
                                                               vector<llf> BM(llf x[N], size_t n){
    ntt(B, len, -1); cls(B, 1, len);
                                                                 size_t f[N]={0},t=0;llf d[N];
                                                                 vector<llf> p[N];
  void pmod(int*A){
                                                                 for(size_t i=1,b=0;i<=n;++i) {</pre>
    static int t[N];
                                                                   for(size_t j=0;j<p[t].size();++j)</pre>
    int l=k+1,len=1;while(len<=(k<<1))len<<=1;</pre>
                                                                     d[i]+=x[i-j-1]*p[t][j];
    cpy(t,A,(k<<1)+1);
                                                                   if(abs(d[i]-=x[i])<=EPS)continue;</pre>
    reverse(t, t+(k<<1)+1);
                                                                   f[t]=i;if(!t){p[++t].resize(i);continue;}
    cls(t,1,len);
                                                                   vector<llf> cur(i-f[b]-1);
    ntt(t,len,1)
                                                                   11f k=-d[i]/d[f[b]];cur.PB(-k);
    for(int i=0;i<len;++i)t[i]=(11)t[i]*ib[i]%mod;</pre>
                                                                   for(size_t j=0;j<p[b].size();j++)</pre>
    ntt(t,len,-1);
                                                                     cur.PB(p[b][j]*k);
    cls(t,1,len)
                                                                   if(cur.size()<p[t].size())cur.resize(p[t].size());</pre>
    reverse(t,t+1);
                                                                   for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];</pre>
    ntt(t,len,1);
                                                                   if(i-f[b]+p[b].size()>=p[t].size()) b=t;
    for(int i=0;i<len;++i)t[i]=(11)t[i]*b[i]%mod;</pre>
                                                                   p[++t]=cur;
    ntt(t,len,-1);
    cls(t,1,len);
                                                                 return p[t];
    for(int i=0;i<k;++i)A[i]=(A[i]-t[i]+mod)%mod;</pre>
                                                               }
    cls(A,k,len);
                                                               5.15 NTT
  void pow(int*A,int n){
    if(n==1) {cls(A,0,k+1);A[1]=1;return;}
    pow(A, n>>1);
                                                               // Remember coefficient are mod P
    int len=1;while(len<=(k<<1))len<<=1;</pre>
                                                               /* p=a*2^n+1
    ntt(A,len,1);
                                                                  n
                                                                       2^n
                                                                                                      root
    for(int i=0;i<len;++i)A[i]=(11)A[i]*A[i]%mod;</pre>
                                                                  16
                                                                       65536
                                                                                     65537
                                                                                                      3
    ntt(A, len, -1);
                                                                       1048576
                                                                                    7340033
                                                                                                      3 */
                                                                  20
    pmod(A);
                                                               // (must be 2<sup>k</sup>)
    if(n&1){
                                                               template<LL P, LL root, int MAXN>
      for(int i=k;i;--i)A[i]=A[i-1];A[0]=0;
                                                               struct NTT{
      pmod(A);
                                                                 static LL bigmod(LL a, LL b) {
                                                                   LL res = 1:
 }
                                                                   for (LL bs = a; b; b >>= 1, bs = (bs * bs) % P)
                                                                     if(b&1) res=(res*bs)%P;
int main(){
                                                                   return res;
 n=rd();k=rd();
  for(int i=1;i<=k;++i)f[i]=(mod+rd())%mod;</pre>
                                                                 static LL inv(LL a, LL b) {
  for(int i=0;i<k;++i)h[i]=(mod+rd())%mod;</pre>
                                                                   if(a==1)return 1:
  for(int i=a[k]=b[k]=1;i<=k;++i)
                                                                   return (((LL)(a-inv(b%a,a))*b+1)/a)%b;
    a[k-i]=b[k-i]=(mod-f[i])%mod;
  int len=1;while(len<=(k<<1))len<<=1;</pre>
                                                                 LL omega[MAXN+1];
```

using VI = vector<int>;

```
NTT() {
    omega[0] = 1;
    LL r = bigmod(root, (P-1)/MAXN);
    for (int i=1; i<=MAXN; i++)</pre>
      omega[i] = (omega[i-1]*r)%P;
  // n must be 2^k
  void tran(int n, LL a[], bool inv_ntt=false){
    int basic = MAXN / n , theta = basic;
    for (int m = n; m >= 2; m >>= 1) {
       int mh = m >> 1;
      for (int i = 0; i < mh; i++) {</pre>
        LL w = omega[i*theta%MAXN];
         for (int j = i; j < n; j += m) {</pre>
           int k = j + mh;
           LL x = a[j] - a[k];
           if (x < 0) x += P;
a[j] += a[k];
           if (a[j] > P) a[j] -= P;
           a[k] = (w * x) % P;
      theta = (theta * 2) % MAXN;
    for (int j = 1; j < n - 1; j++) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
    if (inv_ntt) {
      LL ni = inv(n,P);
      reverse( a+1 , a+n );
for (i = 0; i < n; i++)
        a[i] = (a[i] * ni) % P;
 }
const LL P=2013265921, root=31;
const int MAXN=4194304;
NTT<P, root, MAXN> ntt;
5.16
       Polynomial Operations
```

```
Poly Inverse(Poly f) {
  int n = f.size()
  Poly q(1, fpow(f[0], kMod - 2));
  for (int s = 2;; s <<= 1) {
   if (f.size() < s) f.resize(s);</pre>
    Poly fv(f.begin(), f.begin() + s);
    Poly fq(q.begin(), q.end());
fv.resize(s + s); fq.resize(s + s);
    ntt::Transform(fv, s + s);
    ntt::Transform(fq, s + s);
    for (int i = 0; i < s + s; ++i)
       fv[i] = 1LL * fv[i] * fq[i]%kMod * fq[i]%kMod;
    ntt::InverseTransform(fv, s + s);
    Poly res(s);
    for (int i = 0; i < s; ++i) {
       res[i] = kMod - fv[i];
       if (i < (s >> 1)) {
  int v = 2 * q[i] % kMod;
         (res[i] += v) >= kMod ? res[i] -= kMod : 0;
      }
    }
    q = res;
    if (s >= n) break;
  q.resize(n);
  return q;
Poly Divide(const Poly &a, const Poly &b) {
  int n = a.size(), m = b.size(), k = 2;
  while (k < n - m + 1) k <<= 1;</pre>
  Poly ra(k), rb(k);
  for (int i = 0; i < min(n, k); ++i) ra[i] = a[n-1-i];
for (int i = 0; i < min(m, k); ++i) rb[i] = b[m-1-i];</pre>
  auto rbi = Inverse(rb):
  auto res = Multiply(rbi, ra);
  res.resize(n - m + 1);
  reverse(res.begin(), res.end());
```

```
return res:
Poly Modulo(const Poly &a, const Poly &b) {
  if (a.size() < b.size()) return a;</pre>
  auto dv = Multiply(Divide(a, b), b);
  assert(dv.size() == a.size());
  for (int i = 0; i < dv.size(); ++i)
  dv[i] = (a[i] + kMod - dv[i]) % kMod;</pre>
  while (!dv.empty() && dv.back() == 0) dv.pop_back();
  return dv;
Poly Integral(const Poly &f) {
  int n = f.size();
  VI res(n + 1);
  for (int i = 0; i < n; ++i)</pre>
    res[i+1] = 1LL * f[i] * fpow(i + 1, kMod - 2)%kMod;
  return res:
Poly Evaluate(const Poly &f, const VI &x) {
  if (x.empty()) return Poly();
  int n = x.size();
  vector<Poly> up(n * 2);
  for (int i = 0; i < n; ++i) up[i+n] = {kMod-x[i], 1}; for (int i = n - 1; i > 0; --i)
    up[i] = Multiply(up[i * 2], up[i * 2 + 1]);
  vector<Poly> down(n * 2)
  down[1] = Modulo(f, up[1]);
  for (int i = 2; i < n * 2; ++i)
    down[i] = Modulo(down[i >> 1], up[i]);
  VI y(n);
  for (int i = 0; i < n; ++i) y[i] = down[i + n][0];
  return y;
Poly Interpolate(const VI &x, const VI &y) {
  int n = x.size();
  vector<Poly> up(n * 2);
  for (int i = 0; i < n; ++i) up[i+n] = {kMod-x[i], 1};
  for (int i = n - 1; i > 0; --i)
   up[i] = Multiply(up[i * 2], up[i * 2 + 1]);
  VI a = Evaluate(Derivative(up[1]), x);
  for (int i = 0; i < n; ++i)
a[i] = 1LL * y[i] * fpow(a[i], kMod - 2) % kMod;
  vector<Poly> down(n * 2);
  for (int i = 0; i < n; ++i) down[i + n] = {a[i]};</pre>
  for (int i = n - 1; i > 0; --i)
    auto lhs = Multiply(down[i * 2], up[i * 2 + 1]);
auto rhs = Multiply(down[i * 2 + 1], up[i * 2]);
    assert(lhs.size() == rhs.size());
    down[i].resize(lhs.size());
    for (int j = 0; j < lhs.size(); ++j)
       down[i][j] = (lhs[j] + rhs[j]) % kMod;
  return down[1];
Poly Log(Poly f) {
  int n = f.size();
  if (n == 1) return {0};
  auto d = Derivative(f);
  f.resize(n - 1);
  d = Multiply(d, Inverse(f));
  d.resize(n - 1)
  return Integral(d);
Poly Exp(Poly f) {
  int n = f.size();
  Poly q(1, 1); f[0] += 1;
  for (int s = 1; s < n; s <<= 1) {
     if (f.size() < s + s) f.resize(s + s);</pre>
    Poly g(f.begin(), f.begin() + s + s);
    Poly h(q.begin(), q.end())
    h.resize(s + s); h = Log(h);
for (int i = 0; i < s + s; ++i)
      g[i] = (g[i] + kMod - h[i]) % kMod;
    g = Multiply(g, q);
    g.resize(s + s); q = g;
  assert(q.size() >= n);
  q.resize(n);
  return q:
Poly SquareRootImpl(Poly f) {
 if (f.empty()) return {0};
```

if (t == N) return i

t = (t * B) % P;

11d f = inverse(t, P); for(int i=0;i<=sq+1;i++) {</pre>

if (!R.count(t)) R[t] = i;

```
int z = QuadraticResidue(f[0], kMod), n = f.size();
                                                                  if (R.count(N))
  constexpr int kInv2 = (kMod + 1) >> 1;
                                                                    return i * sq + R[N];
                                                                  N = (N * f) % P;
  if (z == -1) return {-1};
                                                               }
  VI q(1, z);
  for (int s = 1; s < n; s <<= 1) {
                                                                return -1;
    if (f.size() < s + s) f.resize(s + s);</pre>
                                                             }
    VI fq(q.begin(), q.end());
    fq.resize(s + s);
                                                             5.19 Ouadratic residue
    VI f2 = Multiply(fq, fq);
    f2.resize(s + s);
                                                             struct Status{
    for (int i = 0; i < s + s; ++i)
                                                               11 x,y;
      f2[i] = (f2[i] + kMod - f[i]) % kMod;
    f2 = Multiply(f2, Inverse(fq));
                                                             11 w;
    f2.resize(s + s);
                                                             Status mult(const Status& a,const Status& b,ll mod){
    for (int i = 0; i < s + s; ++i)
                                                                Status res
      fq[i] = (fq[i]+kMod - 1LL*f2[i]*kInv2%kMod)%kMod;
                                                                res.x=(a.x*b.x+a.y*b.y%mod*w)%mod;
                                                                res.y=(a.x*b.y+a.y*b.x)%mod;
                                                                return res;
  q.resize(n);
  return q;
                                                             inline Status qpow(Status _base, 11 _pow, 11 _mod) {
                                                                Status res = \{1, 0\};
Poly SquareRoot(Poly f) {
                                                                while(_pow>0){
  int n = f.size(), m = 0;
                                                                  if(_pow&1) res=mult(res,_base,_mod);
  while (m < n \&\& f[m] == 0) m++;
                                                                  _base=mult(_base,_base,_mod);
  if (m == n) return VI(n);
                                                                  _pow>>=1;
  if (m & 1) return {-1}
                                                                }
  auto s = SquareRootImpl(VI(f.begin() + m, f.end()));
                                                                return res;
  if (s[0] == -1) return {-1};
  VI res(n);
                                                             inline 11 check(11 x,11 p){
  for (int i = 0; i < s.size(); ++i) res[i + m/2]=s[i];</pre>
                                                                return qpow_mod(x,(p-1)>>1,p);
  return res:
                                                             inline 11 get_root(11 n,11 p){
                                                                if(p==2) return 1;
5.17 FWT
                                                                if(check(n,p)==p-1) return -1;
                                                                11 a;
/* xor convolution:
                                                                while(true){
 * x = (x0, x1) , y = (y0, y1)
                                                                  a=rand()%p;
 *z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                  w=((a*a-n)%p+p)%p;
                                                                  if(check(w,p)==p-1) break;
 * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
 * z' = ((x\theta+x1)(y\theta+y1)), (x\theta-x1)(y\theta-y1))
* z = (1/2) * z''
                                                                Status res = \{a, 1\}
                                                                res=qpow(res,(p+1)>>1,p);
 * or convolution:
                                                                return res.x;
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
 * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
                                                             5.20 De-Bruijn
const LL MOD = 1e9+7;
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
                                                             int res[maxn], aux[maxn], sz;
  for( int d = 1 ; d < N ; d <<= 1 ) {
                                                             void db(int t, int p, int n, int k) {
  if (t > n) {
    int d2 = d << 1;
    for( int s = 0 ; s < N ; s += d2 )
                                                                  if (n % p == 0)
      for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
  LL ta = x[ i ] , tb = x[ j ];</pre>
                                                                    for (int i = 1; i <= p; ++i)
                                                                     res[sz++] = aux[i];
        x[ i ] = ta+tb;
                                                                } else
        x[ j ] = ta-tb;
                                                                  aux[t] = aux[t - p];
        if( x[ i ] >= MOD ) x[ i ] -= MOD;
                                                                  db(t + 1, p, n, k);
        if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
                                                                  for (int i = aux[t - p] + 1; i < k; ++i) {
                                                                    aux[t] = i;
                                                                    db(t + 1, t, n, k);
  if( inv )
    for( int i = 0 ; i < N ; i++ ) {
  x[ i ] *= inv( N, MOD );</pre>
                                                               }
      x[ i ] %= MOD;
                                                             int de_bruijn(int k, int n) {
                                                                // return cyclic string of len k^n s.t. every string
}
                                                                // of len n using k char appears as a substring.
                                                                if (k == 1) {
5.18 DiscreteLog
                                                                  res[0] = 0;
                                                                  return 1;
// Baby-step Giant-step Algorithm
lld BSGS(lld P, lld B, lld N) {
                                                               for (int i = 0; i < k * n; i++) aux[i] = 0;
  // find B^L = N \mod P
                                                               sz = 0:
  unordered_map<lld, int> R;
                                                               db(1, 1, n, k);
  11d sq = (11d)sqrt(P);
                                                                return sz;
  11d t = 1;
                                                             }
  for (int i = 0; i < sq; i++) {
```

 $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$ and $x_i \ge 0$ for all $1 \le i \le n$. 1. In case of minimization, let $c_i^\prime = -c_i$

Simplex Construction

Standard form: maximize $\sum_{1 \leq i \leq n} c_i x_i$ such that for all $1 \leq j \leq m$,

```
2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
3. \sum_{1 \leq i \leq n} A_{ji} x_i = b_j
           • \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
            • \sum_{1 \le i \le n} A_{ji} x_i \ge b_j
4. If x_i has no lower bound, replace x_i with x_i - x_i'
```

5.22 Simplex

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
  double inv = 1.0 / d[r][s];
  for (int i = 0; i < m + 2; ++i)
    for (int j = 0; j < n + 2; ++j)
  if (i != r && j != s)</pre>
         d[i][j] -= d[r][j] * d[i][s] * inv;
  for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
  d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
  int x = m + z;
  while (true) {
    int s = -1;
     for (int i = 0; i <= n; ++i) {</pre>
       if (!z && q[i] == -1) continue;
if (s == -1 || d[x][i] < d[x][s]) s = i;</pre>
    if (d[x][s] > -eps) return true;
    int r = -1;
for (int i = 0; i < m; ++i) {</pre>
       if (d[i][s] < eps) continue;</pre>
       if (r == -1 ||
         d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
    if (r == -1) return false;
    pivot(r, s);
VD solve(const VVD &a, const VD &b, const VD &c) {
  m = b.size(), n = c.size();
  d = VVD(m + 2, VD(n + 2));
  for (int i = 0; i < m; ++i)
    for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
  p.resize(m), q.resize(n + 1);
  for (int i = 0; i < m; ++i)</pre>
    p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
  for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
  q[n] = -1, d[m + 1][n] = 1;
  int r = 0:
  for (int i = 1; i < m; ++i)</pre>
  if (d[i][n + 1] < d[r][n + 1]) r = i;
if (d[r][n + 1] < -eps) {
    pivot(r, n);
    if (!phase(1) || d[m + 1][n + 1] < -eps)
    return VD(n, -inf);
for (int i = 0; i < m; ++i) if (p[i] == -1) {</pre>
       int s = min_element(d[i].begin(), d[i].end() - 1)
                - d[i].begin();
       pivot(i, s);
    }
  if (!phase(0)) return VD(n, inf);
  VD x(n);
  for (int i = 0; i < m; ++i)
    if (p[i] < n) \times [p[i]] = d[i][n + 1];
  return x;
}}
```

6 Geometry

6.1 Circle Class

```
template<typename T>
struct Circle{
  static constexpr llf EPS = 1e-8;
  Point<T> o; T r;
  vector<Point<llf>> operator&(const Circle& aa)const{
    11f d=o.dis(aa.o);
     if(d>r+aa.r+EPS || d<fabs(r-aa.r)-EPS) return {};
    11f dt = (r*r - aa.r*aa.r)/d, d1 = (d+dt)/2;
    Point<llf> dir = (aa.o-o); dir /= d;
    Point<llf> pcrs = dir*d1 + o;
    dt=sqrt(max(0.0L, r*r - d1*d1)), dir=dir.rot90();
    return {pcrs + dir*dt, pcrs - dir*dt};
};
```

```
6.2 Segment Class
const long double EPS = 1e-8;
template<typename T>
struct Segment{
  // p1.x < p2.x
  Line<T> base;
  Point<T> p1, p2;
  Segment(): base(Line<T>()), p1(Point<T>()), p2(Point<</pre>
      T>()){
    assert(on_line(p1, base) and on_line(p2, base));
                   _, Point<T> __, Point<T> ___): base(_
  Segment(Line<T> _
      ), p1(__), p2(___){
    assert(on_line(p1, base) and on_line(p2, base));
  template<typename T2>
    Segment(const Segment<T2>& _): base(_.base), p1(_.
        p1), p2(_.p2) {}
  typedef Point<long double> Pt;
  friend bool on_segment(const Point<T>& p, const
      Segment& 1){
    if(on_line(p, l.base))
      return (1.p1.x-p.x)*(p.x-1.p2.x)>=0 and (1.p1.y-p
          .y)*(p.y-1.p2.y)>=0;
    return false;
  friend bool have_inter(const Segment& a, const
      Seament& b){
    if(is_parallel(a.base, b.base)){
      return on_segment(a.p1, b) or on_segment(a.p2, b)
           or on_segment(b.p1, a) or on_segment(b.p2, a
    Pt inter = get_inter(a.base, b.base);
    return on_segment(inter, a) and on_segment(inter, b
        );
  friend inline Pt get_inter(const Segment& a, const
      Segment& b){
    if(!have_inter(a, b)){
      return NOT_EXIST;
    }else if(is_parallel(a.base, b.base)){
      if(a.p1 == b.p1){
        if(on_segment(a.p2, b) or on_segment(b.p2, a))
            return INF_P;
        else return a.p1;
      }else if(a.p1 == b.p2){
        if(on_segment(a.p2, b) or on_segment(b.p1, a))
            return INF_P;
        else return a.p1;
      }else if(a.p2 == b.p1){
        if(on_segment(a.p1, b) or on_segment(b.p2, a))
            return INF_P
        else return a.p2;
      }else if(a.p2 == b.p2){
        if(on_segment(a.p1, b) or on_segment(b.p1, a))
            return INF_P
        else return a.p2;
      return INF_P;
```

```
return get_inter(a.base, b.base);
}
friend ostream& operator<<(ostream& ss, const Segment & o){
    ss<<o.base<<", "<<o.p1<<" ~ "<<o.p2;
    return ss;
}
};
template<typename T>
inline Segment<T> get_segment(const Point<T>& a, const Point<T>& b){
    return Segment<T>(get_line(a, b), a, b);
}
```

6.3 Line Class

```
const Point<long double> INF_P(-1e20, 1e20);
const Point<long double> NOT_EXIST(1e20, 1e-20);
template<typename T>
struct Line{
 static constexpr long double EPS = 1e-8;
  // ax+by+c = 0
 T a, b, c;
Line(T _=0, T __=1, T ___=0): a(_), b(__), c(___){
    assert(fabs(a)>EPS or fabs(b)>EPS);}
 template<typename T2>
   Line(const Line<T2>& x): a(x.a), b(x.b), c(x.c){}
  typedef Point<long double> Pt;
 bool equal(const Line& o, true_type) const {
  return fabs(a-o.a)<EPS &&</pre>
    fabs(b-o.b) < EPS && fabs(c-o.b) < EPS;}
 bool equal(const Line& o, false_type) const {
    return a==o.a and b==o.b and c==o.c;}
  bool operator==(const Line& o) const {
    return equal(o, is_floating_point<T>());}
 bool operator!=(const Line& o) const {
    return !(*this == o);}
  friend inline bool on_line__(const Point<T>& p, const
       Line& 1, true_type){
    return fabs(1.a*p.x + 1.b*p.y + 1.c) < EPS;
  friend inline bool on_line__(const Point<T>& p, const
       Line& 1, false_type){
    return 1.a*p.x + 1.b*p.y + 1.c == 0;
 friend inline bool on_line(const Point<T>&p, const
      Line& 1){
    return on_line__(p, 1, is_floating_point<T>());
 friend inline bool is_parallel__(const Line& x, const
       Line& y, true_type){
    return fabs(x.a*y.b - x.b*y.a) < EPS;</pre>
  friend inline bool is_parallel__(const Line& x, const
       Line& y, false_type){
    return x.a*y.b == x.b*y.a;
  friend inline bool is_parallel(const Line& x, const
      Line& y){
    return is_parallel__(x, y, is_floating_point<T>());
 friend inline Pt get_inter(const Line& x, const Line&
       y){
    typedef long double llf;
    if(x==y) return INF_P;
    if(is_parallel(x, y)) return NOT_EXIST;
    llf delta = x.a*y.b - x.b*y.a;
    11f delta_x = x.b*y.c - x.c*y.b;
    11f delta_y = x.c*y.a - x.a*y.c;
    return Pt(delta_x / delta, delta_y / delta);
  friend ostream&operator<<(ostream&ss, const Line&o){</pre>
    ss<<o.a<<"x+"<<o.b<<"y+"<<o.c<<"=0";
    return ss;
 }
template<typename T>
inline Line<T> get_line(const Point<T>& a, const Point<
  return Line<T>(a.y-b.y, b.x-a.x, (b.y-a.y)*a.x-(b.x-a
      .x)*a.y);
```

6.4 Triangle Circumcentre

}

6.5 2D Convex Hull

```
template<typename T>
class ConvexHull_2D{
private:
  typedef Point<T> PT;
  vector<PT> d;
  struct myhash{
    uint64_t operator()(const PT& a) const {
      uint64_t xx=0, yy=0;
      memcpy(&xx, &a.x, sizeof(a.x));
      memcpy(&yy, &a.y, sizeof(a.y));
      uint64_t ret = xx*17+yy*31;
       ret = (ret ^ (ret >> 16))*0x9E3779B1;
      ret = (ret ^ (ret >> 13))*0xC2B2AE35;
      ret = ret ^ xx;
      return (ret ^ (ret << 3)) * yy;</pre>
  };
  unordered_set<PT, myhash> in_hull;
public:
  void init(){in_hull.clear();d.clear();}
  void insert(const PT& x){d.PB(x);}
  void solve(){
    sort(ALL(d), [](const PT& a, const PT& b){
      return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
     vector<PT> s(SZ(d)<<1); int o=0;
     for(auto p: d) {
      while(o \ge 2 \& cross(p-s[o-2], s[o-1]-s[o-2]) <= 0)
        0--
      s[o++] = p;
     for(int i=SZ(d)-2, t = o+1;i>=0;i--){
      while(o>=t\&cross(d[i]-s[o-2],s[o-1]-s[o-2])<=0)
        0-
      s[o++] = d[i];
    s.resize(o-1); swap(s, d);
    for(auto i: s) in_hull.insert(i);
  vector<PT> get(){return d;}
  bool in_it(const PT& x){
     return in_hull.find(x)!=in_hull.end();}
};
```

6.6 2D Farthest Pair

```
// stk is from convex hull
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {
  while(abs(cross(stk[i+1]-stk[i],
        stk[(pos+1)%n]-stk[i])) >
        abs(cross(stk[i+1]-stk[i],
        stk[pos]-stk[i]))) pos = (pos+1)%n;
  ans = max({ans, dis(stk[i], stk[pos]),
        dis(stk[i+1], stk[pos]));
}
```

6.7 2D Closest Pair

```
struct cmp_y {
  bool operator()(const P& p, const P& q) const {
    return p.y < q.y;</pre>
  }
multiset<P, cmp_y> s;
void solve(P a[], int n) {
  sort(a, a + n, [](const P& p, const P& q) {
    return tie(p.x, p.y) < tie(q.x, q.y);</pre>
  11f d = INF; int pt = 0;
  for (int i = 0; i < n; ++i) {
    while (pt < i \text{ and } a[i].x - a[pt].x >= d)
      s.erase(s.find(a[pt++]));
    auto it = s.lower_bound(P(a[i].x, a[i].y - d));
    while (it != s.end() and it->y - a[i].y < d)
      d = min(d, dis(*(it++), a[i]));
    s.insert(a[i]);
  }
}
```

6.8 kD Closest Pair (3D ver.)

```
11f solve(vector<P> v) {
  shuffle(v.begin(), v.end(), mt19937());
  // maybe could replace vector<P> with only P
  unordered_map<lld, unordered_map<lld,
    unordered_map<lld, vector<P>>>> m;
  llf d = dis(v[0], v[1]);
 auto Idx = [&d] (lld x) -> lld {
  return round(x * 2 / d) + 0.1; };
  auto rebuild_m = [&m, &v, &Idx](int k) {
   m.clear();
    for (int i = 0; i < k; ++i)
      m[Idx(v[i].x)][Idx(v[i].y)]
        [Idx(v[i].z)].push_back(v[i]);
  rebuild_m(2);
  for (size_t i = 2; i < v.size(); ++i) {</pre>
    const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
          kz = Idx(v[i].z); bool found = false;
    for (int x = -2; x <= 2; ++x) {
      const 11d nx = x + kx;
      if (m.find(nx) == m.end()) continue;
      auto& mm = m[nx];
      for (int y = -2; y \le 2; ++y) {
        const 11d ny = y + ky;
        if (mm.find(ny) == mm.end()) continue;
        auto& mmm = mm[ny];
        for (int z = -2; z <= 2; ++z) {
          const 11d nz = z + kz;
          if (mmm.find(nz) == mmm.end()) continue;
          for (auto p: mmm[nz]) {
            if (dis(p, v[i]) < d) {</pre>
              d = dis(p, v[i]);
              found = true;
          }
        }
      }
    if (found) rebuild_m(i + 1);
    else m[kx][ky][kz].push_back(v[i]);
  return d:
```

6.9 Simulated Annealing

```
llf anneal() {
  mt19937 rnd_engine( seed );
  uniform_real_distribution< llf > rnd( 0, 1 );
  const llf dT = 0.001;
  // Argument p
  llf S_cur = calc( p ), S_best = S_cur;
  for ( llf T = 2000 ; T > EPS ; T -= dT ) {
    // Modify p to p_prime
    const llf S_prime = calc( p_prime );
    const llf delta_c = S_prime - S_cur;
```

```
llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
if ( rnd( rnd_engine ) <= prob )
    S_cur = S_prime, p = p_prime;
if ( S_prime < S_best ) // find min
    S_best = S_prime, p_best = p_prime;
}
return S_best;
}</pre>
```

6.10 Half Plane Intersection

```
inline int dcmp ( double x ) {
  if( fabs( x ) < eps ) return 0;</pre>
  return x > 0 ? 1 : -1;
struct Line {
  Point st, ed;
  double ang;
  Line(Point _s=Point(), Point _e=Point()):
   st(_s),ed(_e),ang(atan2(_e.y-_s.y,_e.x-_s.x)){}
  inline bool operator< ( const Line& rhs ) const {</pre>
    if(dcmp(ang - rhs.ang) != 0) return ang < rhs.ang;</pre>
    return dcmp( cross( st, ed, rhs.st ) ) < 0;</pre>
  }
};
// cross(pt, line.ed-line.st)>=0 <-> pt in half plane
vector< Line > lns;
deque< Line > que;
deque< Point > pt;
double HPI() {
  sort( lns.begin(), lns.end() );
  que.clear(); pt.clear();
  que.push_back( lns[ 0 ] );
  for ( int i = 1 ; i < (int)lns.size() ; i ++ ) {
   if(!dcmp(lns[i].ang - lns[i-1].ang)) continue;</pre>
    while ( pt.size() > 0 &&
     dcmp(cross(lns[i].st,lns[i].ed,pt.back()))<0){</pre>
      pt.pop_back();que.pop_back();
    while ( pt.size() > 0 &&
     dcmp(cross(lns[i].st,lns[i].ed,pt.front()))<0){</pre>
      pt.pop_front(); que.pop_front();
    pt.push_back(get_point( que.back(), lns[ i ] ));
    que.push_back( lns[ i ] );
  while ( pt.size() > 0 &&
   dcmp(cross(que[0].st, que[0].ed, pt.back()))<0){</pre>
    que.pop_back();
    pt.pop_back();
  while ( pt.size() > 0 &&
   dcmp(cross(que.back().st,que.back().ed,pt[0]))<0){</pre>
    que.pop front():
    pt.pop_front();
  pt.push_back(get_point(que.front(), que.back()));
  vector< Point > conv;
  for ( int i = 0 ; i < (int)pt.size() ; i ++ )</pre>
    conv.push_back( pt[ i ] );
  double ret = 0;
  for ( int i = 1 ; i + 1 < (int)conv.size() ; i ++ )</pre>
    ret += abs(cross(conv[0], conv[i], conv[i + 1]));
  return ret / 2.0;
```

6.11 Ternary Search on Integer

```
int TernarySearch(int 1, int r) {
   // max value @ (1, r]
   while (r - 1 > 1){
      int m = (1 + r)>>1;
      if (f(m) > f(m + 1)) r = m;
      else 1 = m;
   }
   return l+1;
}
```

6.12 Minimum Covering Circle

```
template<typename T>
Circle<llf> MinCircleCover(const vector<PT>& pts){
  random_shuffle(ALL(pts));
  Circle<llf> c = \{pts[0], 0\};
  for(int i=0;i<SZ(pts);i++){</pre>
    if(pts[i].in(c)) continue;
    c = {pts[i], 0};
    for(int j=0;j<i;j++){</pre>
      if(pts[j].in(c)) continue;
      c.o = (pts[i] + pts[j]) / 2;
      c.r = pts[i].dis(c.o);
      for(int k=0;k<j;k++){</pre>
        if(pts[k].in(c)) continue;
        c = get_circum(pts[i], pts[j], pts[k]);
   }
  return c;
```

6.13 KDTree (Nearest Point)

```
const int MXN = 100005;
struct KDTree {
  struct Node {
    int x,y,x1,y1,x2,y2;
    int id,f;
Node *L, *R;
  } tree[MXN], *root;
  int n;
 LL dis2(int x1, int y1, int x2, int y2) {
    LL dx = x1-x2, dy = y1-y2;
    return dx*dx+dy*dy;
  static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
  static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
  void init(vector<pair<int,int>> ip) {
    n = ip.size();
    for (int i=0; i<n; i++) {</pre>
      tree[i].id = i;
      tree[i].x = ip[i].first;
      tree[i].y = ip[i].second;
    root = build_tree(0, n-1, 0);
 Node* build_tree(int L, int R, int d) {
    if (L>R) return nullptr
    int M = (L+R)/2; tree[M].f = d%2;
    nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
    tree[M].x1 = tree[M].x2 = tree[M].x;
    tree[M].y1 = tree[M].y2 = tree[M].y;
    tree[M].L = build_tree(L, M-1, d+1);
    if (tree[M].L) {
      tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
    tree[M].R = build_tree(M+1, R, d+1);
    if (tree[M].R) {
      tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
      tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
    return tree+M;
  int touch(Node* r, int x, int y, LL d2){
    LL dis = sqrt(d2)+1;
    if (x<r->x1-dis || x>r->x2+dis ||
        y<r->y1-dis || y>r->y2+dis)
       return 0;
    return 1:
  void nearest(Node* r,int x,int y,int &mID,LL &md2) {
    if (!r || !touch(r, x, y, md2)) return;
    LL d2 = dis2(r->x, r->y, x, y);
    if (d2 < md2 \mid | (d2 == md2 && mID < r->id)) {
      mID = r->id;
```

```
md2 = d2;
    // search order depends on split dim
    if ((r->f == 0 && x < r->x) ||
        (r->f == 1 && y < r->y))
      nearest(r->L, x, y, mID, md2);
      nearest(r->R, x, y, mID, md2);
    } else {
      nearest(r->R, x, y, mID, md2);
      nearest(r->L, x, y, mID, md2);
  }
  int query(int x, int y) {
    int id = 1029384756;
    LL d2 = 102938475612345678LL;
    nearest(root, x, y, id, d2);
    return id;
} tree;
```

7 Stringology

7.1 Hash

```
class Hash{
private:
    const int p = 127, q = 1051762951;
    int sz, prefix[N], power[N];
    int add(int x, int y){return x+y>=q?x+y-q:x+y;}
    int sub(int x, int y){return x-y<0?x-y+q:x-y;}
    int mul(int x, int y){return 1LL*x*y%q;}
public:
    void init(const string &x){
        sz = x.size();prefix[0]=0;power[0]=1;
        for(int i=1;i<=sz;i++)
            prefix[i]=add(mul(prefix[i-1], p), x[i-1]);
        for(int i=1;i<=sz;i++)power[i]=mul(power[i-1], p);
}
int query(int l, int r){
        // 1-base (l, r]
        return sub(prefix[r], mul(prefix[l], power[r-l]));
}
};</pre>
```

7.2 Suffix Array

```
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
int x[maxn], p[maxn], q[maxn * 2];
// sa[i]: sa[i]-th suffix is the \
// i-th lexigraphically smallest suffix.
// hi[i]: longest common prefix
// of suffix sa[i] and suffix sa[i - 1].
void pre(int *sa, int *c, int n, int z) {
  memset(sa, 0, sizeof(int) * n);
  memcpy(x, c, sizeof(int) * z);
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
  memcpy(x + 1, c, sizeof(int) * (z - 1));
  for (int i = 0; i < n; ++i)
    if (sa[i] && !t[sa[i] - 1])
      sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  memcpy(x, c, sizeof(int) * z);
for (int i = n - 1; i >= 0; --i)
  if (sa[i] && t[sa[i] - 1])
       sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q,
 bool *t, int *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
  memset(c, 0, sizeof(int) * z)
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
  for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
    return;
```

```
for (int i = n - 2; i \ge 0; --i)
    t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  pre(sa, c, n, z);
  for (int i = 1; i <= n - 1; ++i)
if (t[i] && !t[i - 1])
       sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) {
     if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
    bool neq = last < 0 || \</pre>
      memcmp(s + sa[i], s + last,
  (p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
    ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
  pre(sa, c, n, z);
for (int i = nn - 1; i >= 0; --i)
     sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
void build(const string &s) {
  for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];
  _s[(int)s.size()] = 0; // s shouldn't contain 0
  sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];</pre>
  for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
  int ind = 0; hi[0] = 0;
  for (int i = 0; i < (int)s.size(); ++i) {</pre>
     if (!rev[i]) {
       ind = 0:
       continue:
    while (i + ind < (int)s.size() && \</pre>
      s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
    hi[rev[i]] = ind ? ind-- : 0;
}}
```

7.3 Aho-Corasick Algorithm

```
class AhoCorasick{
 private:
    static constexpr int Z = 26;
    struct node{
      node *nxt[ Z ], *fail;
      vector< int > data:
      node(): fail( nullptr ) {
        memset( nxt, 0, sizeof( nxt ) );
        data.clear();
    } *rt;
    inline int Idx( char c ) { return c - 'a'; }
    void init() { rt = new node(); }
    void add( const string& s, int d ) {
      node* cur = rt;
      for ( auto c : s )
        if ( not cur->nxt[ Idx( c ) ] )
          cur->nxt[ Idx( c ) ] = new node();
        cur = cur->nxt[ Idx( c ) ];
      cur->data.push_back( d );
    void compile() {
      vector< node* > bfs;
      size_t ptr = 0;
      for ( int i = 0 ; i < Z ; ++ i ) {
  if ( not rt->nxt[ i ] ) {
          // uncomment 2 lines to make it DFA
          // rt->nxt[i] = rt;
          continue;
        rt->nxt[ i ]->fail = rt;
        bfs.push_back( rt->nxt[ i ] );
      while ( ptr < bfs.size() ) {</pre>
        node* u = bfs[ ptr ++ ];
for ( int i = 0 ; i < Z ; ++ i ) {
           if ( not u->nxt[ i ] ) {
             // u->nxt[i] = u->fail->nxt[i];
             continue;
```

```
node* u_f = u->fail;
          while ( u_f ) {
            if ( not u_f->nxt[ i ] ) {
              u_f = u_f->fail; continue;
            u->nxt[ i ]->fail = u_f->nxt[ i ];
            break;
          if ( not u_f ) u->nxt[ i ]->fail = rt;
          bfs.push_back( u->nxt[ i ] );
        }
     }
    void match( const string& s, vector< int >& ret ) {
      node* u = rt;
      for ( auto c : s ) {
        while ( u != rt and not u->nxt[ Idx( c ) ] )
          u = u->fail;
        u = u->nxt[ Idx( c ) ];
        if ( not u ) u = rt;
        node* tmp = u;
        while ( tmp != rt ) {
          for ( auto d : tmp->data )
            ret.push_back( d );
          tmp = tmp->fail:
      }
    }
} ac;
```

7.4 Suffix Automaton

```
struct Node{
  Node *green, *edge[26];
  int max_len;
  Node(const int _max_len)
    : green(NULL), max_len(_max_len){
    memset(edge, 0, sizeof(edge));
} *ROOT, *LAST;
void Extend(const int c) {
  Node *cursor = LAST
  LAST = new Node((LAST->max_len) + 1);
  for(;cursor&&!cursor->edge[c]; cursor=cursor->green)
    cursor->edge[c] = LAST;
  if (!cursor)
    LAST->green = ROOT;
  else {
    Node *potential_green = cursor->edge[c];
    if((potential_green->max_len)==(cursor->max_len+1))
      LAST->green = potential_green;
//assert(potential_green->max_len>(cursor->max_len+1));
      Node *wish = new Node((cursor->max_len) + 1);
      for(;cursor && cursor->edge[c]==potential_green;
           cursor = cursor->green)
        cursor->edge[c] = wish;
      for (int i = 0; i < 26; i++)
        wish->edge[i] = potential_green->edge[i];
      wish->green = potential_green->green;
      potential_green->green = wish;
      LAST->green = wish;
  }
}
char S[10000001], A[10000001];
int N;
int main(){
  scanf("%d%s", &N, S);
  ROOT = LAST = new Node(0);
  for (int i = 0; S[i]; i++)
    Extend(S[i] - 'a');
  while (N--){
    scanf("%s", A);
    Node *cursor = ROOT;
    bool ans = true;
    for (int i = 0; A[i]; i++){
      cursor = cursor->edge[A[i] - 'a'];
      if (!cursor) {
        ans = false;
```

s += s; int i=0, j=1;

```
National Taiwan University - kiseki
                                                                while (i<n && j<n){</pre>
        break:
                                                                   int k = 0;
      }
                                                                   while (k < n \&\& s[i+k] == s[j+k]) k++;
    puts(ans ? "Yes" : "No");
                                                                   if (s[i+k] \le s[j+k]) j += k+1;
                                                                   else i += k+1;
  return 0;
                                                                   if (i == j) j++;
                                                                int ans = i < n ? i : j;</pre>
                                                                return s.substr(ans, n);
7.5 KMP
vector<int> kmp(const string &s) {
                                                              7.9
                                                                    BWT
  vector<int> f(s.size(), 0);
  /* f[i] = length of the longest prefix
                                                              struct BurrowsWheeler{
     (excluding s[0:i]) such that it coincides
                                                              #define SIGMA 26
     with the suffix of s[0:i] of the same length */
                                                              #define BASE 'a'
  /*i + 1 - f[i] is the length of the
                                                                vector<int> v[ SIGMA ];
     smallest recurring period of s[0:i] */
                                                                void BWT(char* ori, char* res){
  int k = 0;
                                                                   // make ori -> ori + ori
  for (int i = 1; i < (int)s.size(); ++i) {</pre>
                                                                   // then build suffix array
    while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
    if (s[i] == s[k]) ++k;
                                                                void iBWT(char* ori, char* res){
    f[i] = k;
                                                                   for( int i = 0 ; i < SIGMA ; i ++ )</pre>
                                                                    v[ i ].clear();
  return f;
                                                                   int len = strlen( ori );
                                                                   for( int i = 0 ; i < len ; i ++ )</pre>
vector<int> search(const string &s, const string &t) {
                                                                     v[ ori[i] - BASE ].push_back( i );
  // return 0-indexed occurrence of t in s
  vector<int> f = kmp(t), r;
for (int i = 0, k = 0; i < (int)s.size(); ++i)</pre>
                                                                   vector<int> a;
                                                                  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )
for( auto j : v[ i ] ){</pre>
    while(k > 0 && (k==(int)t.size() \mid \mid s[i]!=t[k]))
                                                                       a.push_back( j );
      k = f[k - 1];
                                                                       ori[ ptr ++ ] = BASE + i;
    if (s[i] == t[k]) ++k;
    if (k == (int)t.size()) r.push_back(i-t.size()+1);
                                                                   for( int i = 0 , ptr = 0 ; i < len ; i ++ ){</pre>
                                                                     res[ i ] = ori[ a[ ptr ] ];
  return res;
                                                                     ptr = a[ ptr ];
                                                                   res[ len ] = 0;
7.6 Z value
                                                              } bwt;
char s[MAXN];
int len,z[MAXN];
                                                              7.10 Palindromic Tree
void Z_value() {
  int i,j,left,right;
                                                              struct palindromic_tree{
  z[left=right=0]=len;
                                                                struct node{
  for(i=1;i<len;i++)</pre>
                                                                  int next[26],f,len;
    j=max(min(z[i-left],right-i),0);
                                                                   int cnt, num, st, ed;
    for(;i+j<len&&s[i+j]==s[j];j++);</pre>
                                                                  node(int l=0):f(0),len(1),cnt(0),num(0) {
    if(i+(z[i] = j)>right) {
                                                                     memset(next, 0, sizeof(next)); }
      right=i+z[i];
      left=i;
                                                                vector<node> st;
                                                                vector<char> s;
  }
                                                                int last,n;
}
                                                                void init(){
                                                                   st.clear();s.clear();last=1; n=0;
                                                                   st.push_back(0);st.push_back(-1);
      Manacher
                                                                   st[0].f=1;s.push_back(-1); }
                                                                int getFail(int x){
int z[maxn];
int manacher(const string& s) {
  string t = ".";
                                                                  while(s[n-st[x].len-1]!=s[n])x=st[x].f;
                                                                   return x;}
                                                                void add(int c){
  for(char c:s)) t += c, t += '.';
                                                                   s.push_back(c-='a'); ++n;
  int 1 = 0, r = 0, ans = 0;
                                                                   int cur=getFail(last);
  for (int i = 1; i < t.length(); ++i) {</pre>
                                                                   if(!st[cur].next[c]){
    z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
    while (i - z[i] \ge 0 \& i + z[i] < t.length()) {
                                                                     int now=st.size();
                                                                     st.push_back(st[cur].len+2);
      if(t[i - z[i]] == t[i + z[i]]) ++z[i];
                                                                     st[now].f=st[getFail(st[cur].f)].next[c];
      else break;
                                                                     st[cur].next[c]=now;
                                                                     st[now].num=st[st[now].f].num+1;
    if (i + z[i] > r) r = i + z[i], l = i;
                                                                  last=st[cur].next[c];
  for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);
                                                                  ++st[last].cnt;}
  return ans;
                                                                int size(){ return st.size()-2;}
                                                              } pt;
                                                              int main() {
7.8
      Lexico Smallest Rotation
                                                                string s; cin >> s; pt.init();
                                                                for (int i=0; i<SZ(s); i++) {
  int prvsz = pt.size(); pt.add(s[i]);</pre>
string mcp(string s){
  int n = s.length();
                                                                   if (prvsz != pt.size()) {
```

int r = i, l = r - pt.st[pt.last].len + 1;

// pal @ [1,r]: s.substr(1, r-l+1)

8 Misc

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching an C

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\ldots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + d_2 + \ldots + d_n$ is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all $1 \le k \le n$.

8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W\subseteq X, |W|\leq |N_G(W)|$

8.1.7 Euler's planar graph formula

$$V - E + F = C + 1, E \le 3V - 6$$
(?)

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Lucas's theorem

 ${m\choose n}\equiv\prod_{i=0}^k{m_i\choose n_i}\pmod{p}, \text{ where } m=m_kp^k+m_{k-1}p^{k-1}+\cdots+m_1p+m_0,$ and $n=n_kp^k+n_{k-1}p^{k-1}+\cdots+n_1p+n_0.$

8.2 MaximumEmptuRect

```
int max_empty_rect(int n, int m, bool blocked[N][N]) {
   static int mxu[2][N], me=0, he=1, ans=0;
   for (int i=0;i<m;i++) mxu[he][i]=0;
   for (int i=0;i<n;i++) {
      stack<PII,vector<PII>> stk;
      for (int j=0;j<m;++j) {
        if (blocked[i][j]) mxu[me][j]=0;
        else mxu[me][j]=mxu[he][j]+1;
        int la = j;
      while (!stk.empty()&&stk.top().FF>mxu[me][j]) {
        int x1 = i - stk.top().FF, x2 = i;
        int y1 = stk.top().SS, y2 = j;
        la = stk.top().SS; stk.pop();
        ans=max(ans,(x2-x1)*(y2-y1));
      }
      if (stk.empty()||stk.top().FF<mxu[me][j])
        stk.push({mxu[me][j],la});</pre>
```

```
}
while (!stk.empty()) {
   int x1 = i - stk.top().FF, x2 = i;
   int y1 = stk.top().SS-1, y2 = m-1;
   stk.pop(); ans=max(ans,(x2-x1)*(y2-y1));
   }
   swap(me,he);
}
return ans;
}
```

8.3 DP-opt Condition

8.3.1 totally monotone (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

8.3.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

8.4 Convex 1D/1D DP

```
struct segment {
  int i, l, r;
  segment() {}
  segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
  dp[0] = 0:
  deque<segment> dq; dq.push_back(segment(0, 1, n));
  for (int i = 1; i <= n; ++i) {
    dp[i] = f(dq.front().i, i)
    while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
    dq.front().1 = i + 1;
    segment seg = segment(i, i + 1, n);
    while (dq.size() &&
      f(i, dq.back().1) < f(dq.back().i, dq.back().1))
        dq.pop_back();
    if (dq.size()) {
      int d = 1 << 20, c = dq.back().1;
      while (d \gg 1) if (c + d \ll d, back().r)
        if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
      dq.back().r = c; seg.1 = c + 1;
    if (seg.1 <= n) dq.push_back(seg);</pre>
```

8.5 ConvexHull Optimization

```
inline 1ld DivCeil(1ld n, 1ld d) { // ceil(n/d)
  return n / d + (((n < 0) != (d > 0)) \&\& (n % d));
struct Line {
  static bool flag;
  1ld a, b, l, r; // y=ax+b in [l, r)
  11d operator()(11d x) const { return a * x + b; }
  bool operator<(const Line& i) const {</pre>
    return flag ? tie(a, b) < tie(i.a, i.b) : 1 < i.1;</pre>
  1ld operator&(const Line& i) const {
    return DivCeil(b - i.b, i.a - a);
bool Line::flag = true;
class ConvexHullMax {
  set<Line> L:
 public:
  ConvexHullMax() { Line::flag = true; }
  void InsertLine(lld a, lld b) { // add y = ax + b
    Line now = \{a, b, -INF, INF\};
    if (L.empty()) {
      L.insert(now);
      return;
    Line::flag = true;
    auto it = L.lower_bound(now);
    auto prv = it == L.begin() ? it : prev(it);
```

}

```
if (it != L.end() && ((it != L.begin() &&
                                                            int dp[maxn][2],min_dp[2][2],tmp[2][2],tp[2];
      (*it)(it->1) >= now(it->1) &&
                                                            void dfs(int u,int fa){
      (*prv)(prv->r-1) >= now(prv->r-1)) ||
                                                              if(u<=n){
      (it == L.begin() && it->a == now.a))) return;
                                                                for(int i=0;i<(int)g[u].size();i++){</pre>
    if (it != L.begin())
                                                                  int v=g[u][i];
      while (prv != L.begin() &&
                                                                  if(v==fa) continue;
                                                                  dfs(v,u)
        (*prv)(prv->1) <= now(prv->1))
          prv = --L.erase(prv);
                                                                  memset(tp,0x8f,sizeof tp);
      if (prv == L.begin() && now.a == prv->a)
                                                                  if(v<=n)
        L.erase(prv);
                                                                    tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
                                                                    tp[1]=max
    if (it != L.end())
                                                                      dp[u][0]+dp[v][0]+1
      while (it != --L.end() &&
                                                                      dp[u][1]+max(dp[v][0],dp[v][1])
        (*it)(it->r) <= now(it->r))
                                                                    );
          it = L.erase(it);
                                                                  }else{
    if (it != L.begin()) {
                                                                    tp[0]=dp[u][0]+dp[v][0];
                                                                    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
      prv = prev(it);
      const_cast<Line*>(&*prv)->r=now.l=((*prv)&now);
                                                                  dp[u][0]=tp[0],dp[u][1]=tp[1];
    if (it != L.end())
                                                                }
      const_cast<Line*>(&*it)->l=now.r=((*it)&now);
                                                              }else{
    L.insert(it, now);
                                                                for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                  int v=g[u][i];
  11d Query(11d a) const { // query max at x=a
                                                                  if(v==fa) continue;
    if (L.empty()) return -INF;
                                                                  dfs(v,u);
    Line::flag = false;
    auto it = --L.upper_bound({0, 0, a, 0});
                                                                min_dp[0][0]=0;
    return (*it)(a);
                                                                min_dp[1][1]=1;
  }
                                                                min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
                                                                for(int i=0;i<(int)g[u].size();i++){</pre>
};
                                                                  int v=g[u][i];
                                                                  if(v==fa) continue;
8.6
      Josephus Problem
                                                                  memset(tmp,0x8f,sizeof tmp);
                                                                  tmp[0][0]=max(
// n people kill m for each turn
                                                                    min_dp[0][0]+max(dp[v][0],dp[v][1]),
int f(int n, int m) {
                                                                    min_dp[0][1]+dp[v][0]
  int s = 0;
  for (int i = 2; i <= n; i++)
                                                                  tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
    s = (s + m) \% i;
                                                                  tmp[1][0]=max
  return s;
                                                                    \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
                                                                    min_dp[1][1]+dp[v][0]
// died at kth
int kth(int n, int m, int k){
                                                                  tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
 if (m == 1) return n-1;
                                                                  memcpy(min_dp, tmp, sizeof tmp);
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
                                                                dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
                                                                dp[u][0]=min_dp[0][0];
                                                              }
8.7 Cactus Matching
                                                            int main(){
                                                              int m,a,b;
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
                                                              scanf("%d%d",&n,&m);
                                                              for(int i=0;i<m;i++){</pre>
void tarjan(int u){
  dfn[u]=low[u]=++dfs_idx;
                                                                scanf("%d%d",&a,&b);
                                                                init_g[a].push_back(b);
  for(int i=0;i<(int)init_g[u].size();i++){</pre>
    int v=init_g[u][i];
                                                                init_g[b].push_back(a);
    if(v==par[u]) continue;
                                                              par[1]=-1;
    if(!dfn[v]){
                                                              tarjan(1);
      par[v]=u;
                                                              dfs(1,-1);
      tarjan(v);
                                                              printf("%d\n", max(dp[1][0], dp[1][1]));
      low[u]=min(low[u],low[v]);
      if(dfn[u]<low[v]){</pre>
                                                              return 0:
                                                            }
        g[u].push_back(v);
        g[v].push_back(u);
                                                            8.8 DLX
    }else{
      low[u]=min(low[u],dfn[v]);
      if(dfn[v]<dfn[u]){</pre>
                                                            struct DLX {
        int temp_v=u;
                                                              const static int maxn=210;
        bcc_id++;
                                                              const static int maxm=210;
                                                              const static int maxnode=210*210;
        while(temp_v!=v){
          g[bcc_id+n].push_back(temp_v);
                                                              int n, m, size, row[maxnode], col[maxnode];
          g[temp_v].push_back(bcc_id+n);
                                                              int U[maxnode], D[maxnode], L[maxnode], R[maxnode];
                                                              int H[maxn], S[maxm], ansd, ans[maxn];
          temp_v=par[temp_v];
                                                              void init(int _n, int _m) {
                                                                n = _n, m = _m;
        g[bcc_id+n].push_back(v);
                                                                for(int i = 0; i <= m; ++i) {</pre>
        g[v].push_back(bcc_id+n);
                                                                  S[i] = 0;
        reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
                                                                  U[i] = D[i] = i;
   }
                                                                  L[i] = i-1, R[i] = i+1;
```

R[L[0] = size = m] = 0;

```
for(int i = 1; i <= n; ++i) H[i] = -1;
  void Link(int r, int c) {
    ++S[col[++size] = c];
    row[size] = r; D[size] = D[c];
    U[D[c]] = size; U[size] = c; D[c] = size;
    if(H[r] < 0) H[r] = L[size] = R[size] = size;</pre>
    else {
      R[size] = R[H[r]];
       L[R[H[r]]] = size;
      L[size] = H[r];
       R[H[r]] = size;
    }
  void remove(int c) {
    L[R[c]] = L[c]; R[L[c]] = R[c];
    for(int i = D[c]; i != c; i = D[i])
for(int j = R[i]; j != i; j = R[j]) {
   U[D[j]] = U[j];
         D[U[j]] = D[j];
         --S[col[j]];
  void resume(int c) {
    L[R[c]] = c; R[L[c]] = c;
for(int i = U[c]; i != c; i = U[i])
       for(int j = L[i]; j != i; j = L[j]) {
         U[D[j]] = j;
         D[U[j]] = j;
         ++S[col[j]];
    }
  void dance(int d) {
    if(d>=ansd) return;
    if(R[0] == 0) {
      ansd = d;
      return;
    int c = R[0];
for(int i = R[0]; i; i = R[i])
      if(S[i] < S[c]) c = i;
    remove(c);
    for(int i = D[c]; i != c; i = D[i]) {
      ans[d] = row[i];
       for(int j = R[i]; j != i; j = R[j])
         remove(col[j]);
       dance(d+1);
      for(int j = L[i]; j != i; j = L[j])
         resume(col[j]);
    resume(c);
  }
} sol;
8.9 Tree Knapsack
int dp[N][K];PII obj[N];
vector<int> G[N];
void dfs(int u, int mx){
  for(int s: G[u]) {
    if(mx < obj[s].first) continue;</pre>
    for(int i=0;i<=mx-obj[s].FF;i++)</pre>
    dp[s][i] = dp[u][i];
dfs(s, mx - obj[s].first);
    for(int i=obj[s].FF;i<=mx;i++)</pre>
      dp[u][i] = max(dp[u][i],
         dp[s][i - obj[s].FF] + obj[s].SS);
  }
int main(){
  int n, k; cin >> n >> k;
  for(int i=1;i<=n;i++){</pre>
    int p; cin >> p;
    G[p].push_back(i);
    cin >> obj[i].FF >> obj[i].SS;
  dfs(0, k); int ans = 0;
  for(int i=0;i<=k;i++) ans = max(ans, dp[0][i]);
  cout << ans << '\n';
  return 0;
```

8.10 N Queens Problem

```
vector< int > solve( int n ) {
   // no solution when n=2, 3
   vector< int > ret;
  if ( n % 6 == 2 ) {
  for ( int i = 2 ; i <= n ; i += 2 )
    ret.push_back( i );</pre>
      ret.push_back( 3 ); ret.push_back( 1 );
     for ( int i = 7 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
      ret.push_back( 5 );
  } else if ( n % 6 == 3 ) {
for ( int i = 4 ; i <= n ; i += 2 )</pre>
        ret.push_back( i );
     ret.push_back( 2 );
for ( int i = 5 ; i <= n ; i += 2 )
        ret.push_back( i );
      ret.push_back( 1 ); ret.push_back( 3 );
   } else {
  for ( int i = 2 ; i <= n ; i += 2 )</pre>
     ret.push_back( i );
for ( int i = 1 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
   return ret;
}
```