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	5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17	DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Gauss Elimination Miller Rabin	14 15 15 15 15 16 16 16 16	<pre>#ifdef KISEKI #define safe cerr<<pretty_function\< td=""></pretty_function\<></pre>
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	5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22	DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Gauss Elimination Miller Rabin NTT Range Sieve Partition Number Pi Count (Linear Sieve) Pollard Rho	14 15 15 15 16 16 16 16 16 16	<pre>#ifdef KISEKI #define safe cerr<<pretty_function\< td=""></pretty_function\<></pre>
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6	5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.20 5.21 5.22 6.3 6.4 6.5 6.6 6.7	DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Gauss Elimination Miller Rabin NTT Range Sieve Partition Number Pi Count (Linear Sieve) Pollard Rho Pollynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair	14 15 15 15 16 16 16 16 16 16 17 17 18 18 19 19 19 19 19 20 20 20	<pre>#ifdef KISEKI #define safe cerr<<pretty_function\< td=""></pretty_function\<></pre>
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6	5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.20 5.21 5.23 5.24 6.3 6.4 6.5 6.6 6.7 6.8 6.10 6.11 6.12 6.14 6.15	DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Gauss Elimination Miller Rabin NTT Range Sieve Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair kD Closest Pair kD Closest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Circle Class Intersection of Polygon and Circle Point & Hulls Tangent Convex Hulls Tangent	14 15 15 15 16 16 16 16 16 16 17 17 18 18 19 19 19 20 20 20 20 20 20 20 21 21 21 21	<pre>#ifdef KISEKI #define safe cerr<<pretty_function\< td=""></pretty_function\<></pre>
6	5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.20 5.21 5.22 5.23 5.24 6.5 6.6 6.7 6.8 6.9 6.10 6.11 6.12 6.13 6.14 6.15 6.16	DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Gauss Elimination Miller Rabin NTT Range Sieve Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Circle Class Intersection of Polygon and Circle Point & Hulls Tangent Convex Hulls Tangent Tangent line of Two Circle	14 15 15 15 16 16 16 16 16 16 17 17 18 18 19 19 19 20 20 20 20 20 20 21 21 21 22	<pre>#ifdef KISEKI #define safe cerr<<pretty_function\< td=""></pretty_function\<></pre>
6	5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.20 5.21 5.22 5.24 6.3 6.4 6.5 6.6 6.7 6.8 6.14 6.15 6.14 6.15 6.16 6.16 6.17	DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Gauss Elimination Miller Rabin NTT Range Sieve Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Circle Class Intersection of Polygon and Circle Intersection of Polygon and Circle Point & Hulls Tangent Convex Hulls Tangent Convex Hulls Tangent Tangent line of Two Circle Minimum Covering Circle	14 15 15 15 16 16 16 16 16 16 17 17 18 18 19 19 19 20 20 20 20 20 20 21 21 21 22 22 22	<pre>#ifdef KISEKI #define safe cerr<<pretty_function\< td=""></pretty_function\<></pre>
6	5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.20 5.21 5.22 5.23 5.24 Geo 6.1 6.1 6.1 6.1 6.1 6.1 6.1 6.1 6.1 6.1	DiscreteLog Extended Euler ExtendedFloorSum Fast Fourier Transform FloorSum FWT Gauss Elimination Miller Rabin NTT Range Sieve Partition Number Pi Count (Linear Sieve) Pollard Rho Polynomial Operations Quadratic residue Simplex Simplex Simplex Construction metry Basic Geometry Segment & Line Intersection 2D Convex Hull 3D Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski Sum Circle Class Intersection of Polygon and Circle Point & Hulls Tangent Convex Hulls Tangent Tangent line of Two Circle	14 15 15 15 16 16 16 16 16 16 17 17 18 18 19 19 19 20 20 20 20 20 20 20 21 21 21 22 22 22 22 22	<pre>#ifdef KISEKI #define safe cerr<<pretty_function\< td=""></pretty_function\<></pre>

1.4 Pragma Optimization

```
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popent,abm,mmx,avx,tune=native")
1.5 IO Optimization
static inline int gc() {
 constexpr int B = 1 << 20;
 static char buf[B], *p, *q;
 if(p == q \&\&
  (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
  return EOF:
 return *p++;
template < typename T >
static inline bool gn( T &x ) {
 int c = gc(); T sgn = 1; x = 0;
while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
if(c == '-') sgn = -1, c = gc();
 if(c == EOF) return false;
 while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
 return x *= sgn, true;
```

#pragma GCC optimize("Ofast,no-stack-protector")

2 Data Structure

2.1 Dark Magic

```
2.2 Link-Cut Tree
template <typename Val> class LCT {
private:
 struct node {
  int pa, ch[2];
  bool rev;
  \label{eq:valv} \begin{array}{lll} \mbox{Val } \mbox{v}_{-}\mbox{prod}, \ \mbox{v}_{-}\mbox{prod}; \\ \mbox{node}() \mbox{:} \mbox{pa}\{\theta\}, \mbox{ ch}\{\emptyset, \mbox{ } \emptyset\}, \mbox{ rev}\{\mbox{false}\}, \mbox{ } v_{-}\mbox{prod}\{\}, \end{array}
     v_rprod{} {};
 }:
 vector<node> nodes;
 set<pair<int, int>> edges;
 bool is_root(int u) const {
  const int p = nodes[u].pa;
  return nodes[p].ch[0] != u and nodes[p].ch[1] != u;
 bool is_rch(int u) const {
  return (not is_root(u)) and nodes[nodes[u].pa].ch[1]
     == u;
 void down(int u) {
  if (auto &cnode = nodes[u]; cnode.rev) {
   if (cnode.ch[0]) set_rev(cnode.ch[0]);
   if (cnode.ch[1]) set_rev(cnode.ch[1]);
   cnode.rev = false;
  }
 }
 void up(int u) {
  auto &cnode = nodes[u];
  cnode.v_prod =
   nodes[cnode.ch[0]].v_prod * cnode.v * nodes[cnode.ch
     [1]].v_prod;
  cnode.v_rprod =
   nodes[cnode.ch[1]].v_rprod * cnode.v * nodes[cnode.
     ch[0]].v_rprod;
 void set_rev(int u) {
```

```
swap(nodes[u].ch[0], nodes[u].ch[1]);
  swap(nodes[u].v_prod, nodes[u].v_rprod);
  nodes[u].rev ^= 1;
 void rotate(int u) {
  int f = nodes[u].pa, g = nodes[f].pa, l = is_rch(u);
if (nodes[u].ch[l ^ 1])
  nodes[nodes[u].ch[1 ^ 1]].pa = f;
  if (not is_root(f))
  nodes[g].ch[is_rch(f)] = u;
  nodes[f].ch[1] = nodes[u].ch[1 ^ 1];
  nodes[u].ch[1 ^ 1] = f;
 nodes[u].pa = g, nodes[f].pa = u;
  up(f);
 void splay(int u) {
  vector<int> stk = {u};
  while (not is_root(stk.back()))
   stk.push_back(nodes[stk.back()].pa);
  for (; not stk.empty(); stk.pop_back())
  down(stk.back())
  for(int f=nodes[u].pa;!is_root(u);f=nodes[u].pa){
   if(!is_root(f))rotate(is_rch(u)==is_rch(f)?f:u);
  rotate(u):
 up(u);
 void access(int u) {
  int last = 0;
  for (int last = 0; u; last = u, u = nodes[u].pa) {
   splay(u);
   nodes[u].ch[1] = last;
   up(u);
 }
 int find_root(int u) {
 access(u); splay(u);
  int la = 0;
  for (; u; la = u, u = nodes[u].ch[0]) down(u);
  return la:
 void change_root(int u) {
 access(u); splay(u); set_rev(u);
 void link(int x, int y) {
  change_root(y); nodes[y].pa = x;
 void split(int x, int y) {
 change_root(x); access(y); splay(y);
 void cut(int x, int y) {
  split(x, y)
 nodes[y].ch[0] = nodes[x].pa = 0;
  up(y);
public:
 LCT(int n = 0) : nodes(n + 1) {}
 int add(const Val &v = {}) {
 nodes.push_back(v);
  return int(nodes.size()) - 2;
 int add(Val &&v) {
 nodes.emplace_back(move(v));
  return int(nodes.size()) - 2;
 void set_val(int u, const Val &v) {
 splay(++u); nodes[u].v = v; up(u);
 Val query(int x, int y) {
  split(++x, ++y);
  return nodes[y].v_prod;
 bool connected(int u, int v) { return find_root(++u)
    == find_root(++v); }
 void add_edge(int u, int v) {
  if (++u > ++v) swap(u, v);
 edges.emplace(u, v); link(u, v);
 void del_edge(int u, int v) {
  auto k = minmax(++u, ++v);
  if (auto it = edges.find(k); it != edges.end()) {
   edges.erase(it); cut(u, v);
```

```
int getRank(node *o) {
}
                                                             int r = sz(o->lc);
};
                                                             for (;o->pa != nullptr; o = o->pa)
                                                              if (o->pa->rc != o) r += sz(o->pa->lc);
2.3
     LiChao Segment Tree
                                                             return r:
struct L {
int m, k, id;
L() : id(-1) {}
                                                            #undef sz
                                                           }
L(int a, int b, int c) : m(a), k(b), id(c) {}
                                                           2.5 Linear Basis
int at(int x) { return m * x + k; }
                                                           template <int BITS>
                                                           struct LinearBasis {
class LiChao {
private:
                                                            array<uint64_t, BITS> basis;
int n; vector<L> nodes;
                                                            Basis() { basis.fill(0); }
                                                            void add(uint64_t x)
static int lc(int x) { return 2 * x + 1; }
                                                             for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
static int rc(int x) { return 2 * x + 2; }
                                                              if (basis[i] == 0) {
 void insert(int 1, int r, int id, L ln) {
 int m = (1 + r) >> 1;
                                                               basis[i] = x;
 if (nodes[id].id == -1) {
                                                               return;
  nodes[id] = ln;
                                                              x ^= basis[i];
   return;
                                                             }
 bool atLeft = nodes[id].at(1) < ln.at(1);</pre>
 if (nodes[id].at(m) < ln.at(m)) {</pre>
                                                            bool ok(uint64_t x) {
  atLeft ^= 1;
                                                             for (int i = 0; i < BITS; ++i)</pre>
  swap(nodes[id], ln);
                                                              if ((x >> i) & 1) x ^= basis[i];
                                                             return x == 0;
 if (r - l == 1) return;
 if (atLeft) insert(1, m, lc(id), ln);
                                                           };
  else insert(m, r, rc(id), ln);
                                                                Binary Search On Segment Tree
                                                           2.6
                                                           // find_first = x -> minimal x s.t. check( [a, x) )
int query(int 1, int r, int id, int x) {
  int ret = 0, m = (1 + r) >> 1;
                                                           if (nodes[id].id != -1)
                                                           template <typename C>
  ret = nodes[id].at(x);
                                                           int find_first(int 1, const C &check) {
  if (r - 1 == 1) return ret;
                                                            if (1 >= n) return n;
 if (x < m) return max(ret, query(1, m, lc(id), x));</pre>
                                                            1 += sz;
  return max(ret, query(m, r, rc(id), x));
                                                            for (int i = height; i > 0; i--)
                                                             propagate(1 >> i);
                                                            Monoid sum = identity;
public:
while ((1 & 1) == 0) 1 >>= 1;
                                                             if (check(f(sum, data[1]))) {
int query(int x) { return query(0, n, 0, x); }
                                                              while (1 < sz) {</pre>
}:
                                                               propagate(1);
                                                               1 <<= 1:
2.4 Treap
                                                               auto nxt = f(sum, data[1]);
namespace Treap{
                                                               if (not check(nxt)) {
#define sz( x ) ( ( x ) ? ( ( x )->size ) : \theta )
                                                                sum = nxt;
struct node{
                                                                1++;
  int size;
                                                               }
 uint32_t pri;
                                                              }
 node *1c, *rc, *pa;
                                                              return 1 + 1 - sz;
 node():size(0),pri(rand()),lc(0),rc(0),pa(0){}
                                                             }
 void pull() {
                                                             sum = f(sum, data[1++]);
  size = 1; pa = nullptr;
                                                            } while ((1 & -1) != 1);
  if ( lc ) { size += lc->size; lc->pa = this; }
                                                            return n;
   if ( rc ) { size += rc->size; rc->pa = this; }
                                                           template <typename C>
                                                           int find_last(int r, const C &check) {
}:
node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
                                                            if (r <= 0) return -1;
                                                            r += sz;
 if ( L->pri > R->pri ) {
                                                            for (int i = height; i > 0; i--)
  L->rc = merge( L->rc, R ); L->pull();
                                                             propagate((r - 1) >> i);
   return L;
                                                            Monoid sum = identity;
  } else {
                                                            do {
   R->lc = merge( L, R->lc ); R->pull();
                                                             while (r > 1 \text{ and } (r \& 1)) r >>= 1;
   return R;
                                                             if (check(f(data[r], sum))) {
  }
                                                              while (r < sz) {
void split_by_size( node*rt,int k,node*&L,node*&R ) {
  if ( not rt ) L = R = nullptr;
                                                               propagate(r);
                                                               r = (r << 1) + 1;
 else if( sz( rt->lc ) + 1 <= k ) {
                                                               auto nxt = f(data[r], sum);
                                                               if (not check(nxt)) {
  L = rt
   split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                sum = nxt;
   L->pull();
                                                               }
  } else {
                                                              }
  R = rt:
   split_by_size( rt->lc, k, L, R->lc );
                                                              return r - sz;
   R->pull();
                                                             }
                                                             sum = f(data[r], sum);
}
                                                            } while ((r & -r) != r);
```

```
low[u] = min(low[u], low[v]);
return -1;
                                                                if (low[v] > dfn[u]) bridge[t] = true;
3
    Graph
                                                             public:
    2-SAT (SCC)
                                                              void init(int n_) {
                                                               G.clear(); G.resize(n = n_);
class TwoSat{
                                                               low.assign(n, ecnt = 0);
private:
                                                               dfn.assign(n, 0);
int n;
vector<vector<int>> rG,G,sccs;
                                                              void add_edge(int u, int v) {
vector<int> ord,idx;
                                                               G[u].emplace_back(v, ecnt);
vector<bool> vis,result;
                                                               G[v].emplace_back(u, ecnt++);
void dfs(int u){
 vis[u]=true
                                                              void solve() {
 for(int v:G[u])
                                                               bridge.assign(ecnt, false);
  if(!vis[v]) dfs(v);
                                                               for (int i = 0; i < n; ++i)
 ord.push_back(u);
                                                                if (not dfn[i]) dfs(i, i);
void rdfs(int u){
                                                              bool is_bridge(int x) { return bridge[x]; }
 vis[u]=false;idx[u]=sccs.size()-1;
                                                            } bcc_bridge;
 sccs.back().push_back(u);
  for(int v:rG[u])
                                                            3.3 BCC Vertex
  if(vis[v])rdfs(v);
                                                            class BCC_AP {
                                                             private:
public:
                                                              int n, ecnt;
void init(int n_){
                                                              vector<vector<pair<int,int>>> G;
 G.clear();G.resize(n=n_);
                                                              vector<int> bcc, dfn, low, st;
 rG.clear();rG.resize(n);
                                                              vector<bool> ap, ins;
void dfs(int u, int f)
 sccs.clear();ord.clear();
 idx.resize(n);result.resize(n);
                                                               dfn[u] = low[u] = dfn[f] + 1;
                                                               int ch = 0;
void add_edge(int u,int v){
                                                               for (auto [v, t]: G[u]) if (v != f) {
 G[u].push_back(v);rG[v].push_back(u);
                                                                if (not ins[t]) {
                                                                 st.push_back(t);
void orr(int x,int y){
                                                                 ins[t] = true;
 if ((x^y)==1)return;
 add_edge(x^1,y); add_edge(y^1,x);
                                                                if (dfn[v]) {
                                                                 low[u] = min(low[u], dfn[v]);
bool solve(){
                                                                 continue:
 vis.clear();vis.resize(n);
                                                                 } ++ch; dfs(v, u);
 for(int i=0;i<n;++i)</pre>
                                                                low[u] = min(low[u], low[v]);
  if(not vis[i])dfs(i);
                                                                if (low[v] >= dfn[u]) {
  reverse(ord.begin(),ord.end());
                                                                 ap[u] = true;
 for (int u:ord){
                                                                 while (true) {
  if(!vis[u])continue;
                                                                  int eid = st.back(); st.pop_back();
  sccs.push_back(vector<int>());
                                                                  bcc[eid] = ecnt;
  rdfs(u);
                                                                  if (eid == t) break;
 for(int i=0;i<n;i+=2)</pre>
                                                                 ecnt++:
  if(idx[i]==idx[i+1])
                                                                }
    return false
 vector<bool> c(sccs.size());
                                                               if (ch == 1 and u == f) ap[u] = false;
 for(size_t i=0;i<sccs.size();++i){</pre>
   for(auto sij : sccs[i]){
                                                             public:
    result[sij]=c[i]
                                                              void init(int n_) {
    c[idx[sij^1]]=!c[i];
                                                               G.clear(); G.resize(n = n_);
                                                               ecnt = 0; ap.assign(n, false);
                                                               low.assign(n, 0); dfn.assign(n, 0);
 return true;
                                                              void add_edge(int u, int v) {
bool get(int x){return result[x];}
                                                               G[u].emplace_back(v, ecnt)
int get_id(int x){return idx[x];}
                                                               G[v].emplace_back(u, ecnt++);
int count(){return sccs.size();}
                                                              }
} sat2;
                                                              void solve() {
                                                               ins.assign(ecnt, false);
3.2 BCC Edge
                                                               bcc.resize(ecnt); ecnt = 0;
                                                               for (int i = 0; i < n; ++i)
class BCC_Bridge {
                                                                if (not dfn[i]) dfs(i, i);
private:
 int n, ecnt;
 vector<vector<pair<int,int>>> G;
                                                              int get_id(int x) { return bcc[x]; }
 vector<int> dfn, low;
                                                              int count() { return ecnt; }
 vector<bool> bridge;
                                                              bool is_ap(int x) { return ap[x]; }
 void dfs(int u, int f) {
  dfn[u] = low[u] = dfn[f] + 1;
                                                            } bcc_ap;
                                                            3.4 Centroid Decomposition
  for (auto [v, t]: G[u]) {
   if (v == f) continue;
if (dfn[v]) {
                                                            struct Centroid {
                                                             vector<vector<int64_t>> Dist;
    low[u] = min(low[u], dfn[v]);
                                                             vector<int> Parent, Depth;
                                                             vector<int64_t> Sub, Sub2;
     continue;
                                                             vector<int> Sz, Sz2;
    dfs(v, u);
                                                             Centroid(vector<vector<pair<int, int>>> g) {
```

```
int N = g.size();
 vector<bool> Vis(N);
 vector<int> sz(N), mx(N);
 vector<int> Path;
Dist.resize(N);
Parent.resize(N);
 Depth.resize(N)
 auto DfsSz = [&](auto dfs, int x) -> void {
  Vis[x] = true; sz[x] = 1; mx[x] = 0;
  for (auto [u, w] : g[x]) {
   if (Vis[u]) continue;
   dfs(dfs, u);
   sz[x] += sz[u]
   mx[x] = max(mx[x], sz[u]);
  Path.push_back(x);
 auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
  -> void {
  Dist[x].push_back(D);Vis[x] = true;
 for (auto [u, w] : g[x]) {
   if (Vis[u]) continue;
   dfs(dfs, u, D + w);
  }
 auto Dfs = [&]
  (auto dfs, int x, int D = 0, int p = -1)->void {
  Path.clear(); DfsSz(DfsSz, x);
  int M = Path.size();
  int C = -1;
  for (int u : Path) {
   if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
   Vis[u] = false;
  DfsDist(DfsDist, C);
  for (int u : Path) Vis[u] = false;
  Parent[C] = p; Vis[C] = true;
  Depth[C] = D;
  for (auto [u, w] : g[C]) {
  if (Vis[u]) continue
   dfs(dfs, u, D + 1, C);
 Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
Sz.resize(N); Sz2.resize(N);
void Mark(int v) {
int x = v, z = -1
 for (int i = Depth[v]; i >= 0; --i) {
 Sub[x] += Dist[v][i]; Sz[x]++;
 if (z != -1) {
   Sub2[z] += Dist[v][i];
   Sz2[z]++;
  z = x; x = Parent[x];
}
int64_t Query(int v) {
int64_t res = 0;
 int x = v, z = -1;
for (int i = Depth[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
 if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
 z = x; x = Parent[x];
return res;
}
    Directed Minimum Spanning Tree
```

```
struct DirectedMST { // find maximum
    struct Edge {
    int u, v;
    int w;
    Edge(int u, int v, int w) : u(u), v(v), w(w) {}
};
    vector<Edge> Edges;
    void clear() { Edges.clear(); }
    void addEdge(int a, int b, int w) { Edges.emplace_back
        (a, b, w); }
    int solve(int root, int n) {
        vector<Edge> E = Edges;
    }
}
```

fa[x] = p

return c ? p : val[x];

vector<int> build(int s, int n) {

```
int ans = 0;
  while (true) {
   // find best in edge
   vector<int> in(n, -inf), prv(n, -1);
   for (auto e : È)
    if (e.u != e.v && e.w > in[e.v]) {
     in[e.v] = e.w;
     prv[e.v] = e.u;
   in[root] = 0;
   prv[root] = -1;
   for (int i = 0; i < n; i++)
    if (in[i] == -inf)
     return -inf;
   // find cycle
   int tot = 0:
   for (int i = 0; i < n; i++) {
    ans += in[i];
    for (int x = i; x != -1 && id[x] == -1; x = prv[x])
     if (vis[x] == i) {
      for (int y = prv[x]; y != x; y = prv[y])
id[y] = tot;
      id[x] = tot++;
      break;
     vis[x] = i;
   if (!tot)
    return ans;
   for (int i = 0; i < n; i++)</pre>
    if (id[i] == -1)
     id[i] = tot++;
   for (auto &e : E) {
    if (id[e.u] != id[e.v])
     e.w -= in[e.v];
    e.u = id[e.u], e.v = id[e.v];
   n = tot;
   root = id[root];
  assert(false);
} DMST;
3.6 Dominator Tree
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
int dom[maxn], val[maxn], rp[maxn], tk;
void init(int n) {
 // vertices are numbered from 0 to n-1
 fill(dfn, dfn + n, -1); fill(rev, rev + n, -1);
 fill(fa, fa + n, -1); fill(val, val + n, -1);
 fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
 fill(dom, dom + n, -1); tk = 0;
 for (int i = 0; i < n; ++i) {
  g[i].clear(); r[i].clear(); rdom[i].clear();
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
 rev[dfn[x] = tk] = x;
 fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
 for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
  r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
   if (fa[x] == x) return c ? -1 : x;
 int p = find(fa[x], 1);
 if (p == -1) return c ? fa[x] : val[x];
if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
```

// head[i] : head of the chain i

```
// return the father of each node in the dominator tree
                                                                // chian[u] : chain id of the chain u is on
                                                                void predfs(int u, int f) {
  dep[u] = dep[pa[u] = f] + 1;
// p[i] = -2 if i is unreachable from s
 dfs(s);
                                                                 for (int v : G[u]) if (v != f) {
 for (int i = tk - 1; i >= 0; --i) {
 for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
                                                                  predfs(v, u);
  if (i) rdom[sdom[i]].push_back(i);
                                                                  if (lowbit(chain[u]) < lowbit(chain[v]))</pre>
  for (int &u : rdom[i]) {
                                                                   chain[u] = chain[v];
   int p = find(u);
   if (sdom[p] == i) dom[u] = i;
                                                                 if (chain[u] == 0) chain[u] = ++chains;
   else dom[u] = p;
                                                                void dfschain(int u, int f) {
  if (i) merge(i, rp[i]);
                                                                 tl[u] = timer++;
 }
                                                                 if (head[chain[u]] == -1)
 vector<int> p(n, -2); p[s] = -1;
                                                                  head[chain[u]] = u;
 for (int i = 1; i < tk; ++i)</pre>
                                                                 for (int v : G[u])
 if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
                                                                  if (v != f and chain[v] == chain[u])
                                                                 dfschain(v, u);
for (int v : G[u])
 for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
 return p;
                                                                  if (v != f and chain[v] != chain[u])
}}
                                                                   dfschain(v, u);
3.7 Edge Coloring
                                                                 tr[u] = timer;
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
                                                               public:
                                                                LBD(int n) : timer(0), chains(0), G(n), tl(n), tr(n),
                                                                chain(n), head(n, -1), dep(n), pa(n) {}
void add_edge(int u, int v) {
 for (int i = 0; i <= N; i++)
  for (int j = 0; j <= N; j++)
C[i][j] = G[i][j] = 0;
                                                                 G[u].push_back(v); G[v].push_back(u);
                                                                void decompose() { predfs(0, 0); dfschain(0, 0); }
void solve(vector<pair<int, int>> &E, int N) {
 int X[kN] = {}, a;
auto update = [&](int u) {
                                                                PII get_subtree(int u) { return {tl[u], tr[u]}; }
                                                                vector<PII> get_path(int u, int v) {
                                                                 vector<PII> res;
while (chain[u] != chain[v]) +
  for (X[u] = 1; C[u][X[u]]; X[u]++);
                                                                  if (dep[head[chain[u]]] < dep[head[chain[v]]])</pre>
 auto color = [&](int u, int v, int c) {
 int p = G[u][v];
G[u][v] = G[v][u] = c;
                                                                   swap(u, v)
                                                                  int s = head[chain[u]];
  C[u][c] = v, C[v][c] = u;
                                                                  res.emplace_back(tl[s], tl[u] + 1);
  C[u][p] = C[v][p] = 0;
                                                                  u = pa[s];
  if (p) X[u] = X[v] = p
  else update(u), update(v);
                                                                 if (dep[u] < dep[v]) swap(u, v);</pre>
  return p;
                                                                 res.emplace_back(tl[v], tl[u] + 1);
                                                                 return res;
 }:
 auto flip = [&](int u, int c1, int c2) {
                                                                }
 int p = C[u][c1];
                                                              };
  swap(C[u][c1], C[u][c2]);
                                                                     Manhattan Minimum Spanning Tree
                                                               3.9
  if (p) G[u][p] = G[p][u] = c2;
  if (!C[u][c1]) X[u] = c1;
                                                               typedef Point<int> P;
                                                               vector<array<int, 3>> manhattanMST(vector<P> ps) {
  if (!C[u][c2]) X[u] = c2;
                                                                vi id(sz(ps));
                                                                iota(all(id), 0);
 }:
 for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {
                                                                vector<array<int, 3>> edges;
                                                                rep(k, 0, 4) {
                                                                 sort(all(id),
 auto [u, v] = E[t];
                                                                                [&](int i, int j) {
  int v0 = v, c = X[u], c0 = c, d;
                                                                  return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;</pre>
  vector<pair<int, int>> L; int vst[kN] = {};
                                                                 });
  while (!G[u][v0]) {
                                                                 map<int, int> sweep;
   L.emplace_back(v, d = X[v]);
                                                                 for (int i : id) {
   if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
                                                                  for (auto it = sweep.lower_bound(-ps[i].y);
     c = color(u, L[a].first, c);
                                                                     it != sweep.end(); sweep.erase(it++)) {
   else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
                                                                   int j = it->second;
                                                                   P d = ps[i] - ps[j];
     color(u, L[a].first, L[a].second);
   else if (vst[d]) break
                                                                   if (d.y > d.x) break;
   else vst[d] = 1, v = C[u][d];
                                                                   edges.push_back({d.y + d.x, i, j});
  if (!G[u][v0]) {
                                                                  sweep[-ps[i].y] = i;
   for (; v; v = flip(v, c, d), swap(c, d));
if (C[u][c0]) { a = int(L.size()) - 1;
                                                                 for (P &p : ps)
    while (--a >= 0 && L[a].second != c);
                                                                  if (k \& 1) p.x = -p.x;
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
                                                                  else swap(p.x, p.y);
   } else t--;
                                                                return edges; // [{w, i, j}, ...]
                                                               }
                                                               3.10 MaxClique
3.8 Lowbit Decomposition
                                                               // contain a self loop u to u, than u won't in clique
class LBD {
                                                               template < size_t MAXN >
                                                               class MaxClique{
 int timer, chains;
 vector<vector<int>> G;
                                                               private:
 vector<int> tl, tr, chain, head, dep, pa;
                                                                using bits = bitset< MAXN >;
 // chains : number of chain
                                                                bits popped, G[ MAXN ], ans
 // tl, tr[u] : subtree interval in the seq. of u
                                                                size_t deg[ MAXN ], deo[ MAXN ], n;
```

void sort_by_degree() {

cs[k][p] = 1;

```
popped.reset();
                                                                       if (k < km) r[t++] = p;
  for ( size_t i = 0 ; i < n ; ++ i )
    deg[ i ] = G[ i ].count();
                                                                      c.resize(m);
  for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t mi = MAXN, id = 0;</pre>
                                                                      if (t) c[t - 1] = 0;
                                                                      for (int k = km; k <= mx; k++) {</pre>
    for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )
    mi = deg[ id = j ];</pre>
                                                                       for (int p = int(cs[k]._Find_first());
                                                                           p < kN; p = int(cs[k]._Find_next(p))) {
                                                                        r[t] = p; c[t++] = k;
    popped[ deo[ i ] = id ] = 1;
    for( size_t u = G[ i ]._Find_first() ;
     u < n ; u = G[ i ]._Find_next( u ) )</pre>
                                                                     void dfs(vector<int> &r, vector<int> &c, int 1,
       -- deg[ u ];
  }
                                                                      bitset<kN> mask) {
}
                                                                      while (!r.empty()) {
                                                                       int p = r.back(); r.pop_back();
void BK( bits R, bits P, bits X ) {
  if (R.count()+P.count() <= ans.count()) return;</pre>
                                                                       mask[p] = 0;
  if ( not P.count() and not X.count() )
                                                                       if (q + c.back() <= ans) return;</pre>
  if ( R.count() > ans.count() ) ans = R;
                                                                       cur[q++] = p;
                                                                       vector<int> nr, nc;
   return;
                                                                       bitset<kN> nmask = mask & a[p];
  /* greedily chosse max degree as pivot
                                                                       for (int i : r)
  bits cur = P | X; size_t pivot = 0, sz = 0;
                                                                        if (a[p][i]) nr.push_back(i);
  for ( size_t u = cur._Find_first() ;
                                                                       if (!nr.empty()) {
                                                                        if (1 < 4) {
  u < n ; u = cur._Find_next( u ) )</pre>
    if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
                                                                         for (int i : nr)
  cur = P & ( ~G[ pivot ] );
                                                                           d[i] = int((a[i] & nmask).count());
  */ // or simply choose first
                                                                          sort(nr.begin(), nr.end(),
  bits cur = P & (~G[ ( P | X )._Find_first() ]);
                                                                           [&](int x, int y)
                                                                            return d[x] > d[y];
  for ( size_t u = cur._Find_first()
   u < n ; u = cur._Find_next( u ) ) {
   if ( R[ u ] ) continue;
                                                                       csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
   R[u] = 1;
   BK( R, P & G[ u ], X & G[ u ] )
   R[u] = P[u] = 0, X[u] = 1;
                                                                        ans = q; copy(cur, cur + q, sol);
                                                                       c.pop_back(); q--;
public:
 void init( size_t n_ ) {
  n = n_{-};
                                                                     int solve(bitset<kN> mask) { // vertex mask
  for ( size_t i = 0 ; i < n ; ++ i )
                                                                      vector<int> r, c;
                                                                      for (int i = 0; i < n; i++)
  if (mask[i]) r.push_back(i);</pre>
   G[ i ].reset();
  ans.reset();
                                                                      for (int i = 0; i < n; i++)
void add_edges( int u, bits S ) { G[ u ] = S; }
                                                                       d[i] = int((a[i] & mask).count());
void add_edge( int u, int v ) {
  G[ u ][ v ] = G[ v ][ u ] = 1;
                                                                      sort(r.begin(), r.end(),
  [&](int i, int j) { return d[i] > d[j]; });
                                                                      csort(r, c);
                                                                      dfs(r, c, 1, mask);
return ans; // sol[0 ~ ans-1]
int solve() {
  sort_by_degree(); // or simply iota( deo... )
  for ( size_t i = 0 ; i < n ; ++ i )
   deg[ i ] = G[ i ].count();
                                                                   } graph;
  bits pob, nob = 0; pob.set();
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;
                                                                    3.12 Minimum Mean Cycle
  for ( size_t i = 0 ; i < n ; ++ i ) {</pre>
                                                                   /* minimum mean cycle O(VE) */
   size_t v = deo[ i ];
                                                                    struct MMC{
   bits tmp; tmp[ v ] = 1;
                                                                   #define FZ(n) memset((n),0,sizeof(n))
   BK( tmp, pob & G[ v ], nob & G[ v ] );
pob[ v ] = 0, nob[ v ] = 1;
                                                                    #define E 101010
                                                                    #define V 1021
                                                                    #define inf 1e9
  return static_cast< int >( ans.count() );
                                                                     struct Edge { int v,u; double c; };
                                                                     int n, m, prv[V][V], prve[V][V], vst[V];
};
                                                                     Edge e[E];
                                                                     vector<int> edgeID, cycle, rho;
3.11 MaxCliqueDyn
                                                                     double d[V][V];
                                                                     void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
constexpr int kN = 150;
struct MaxClique { // Maximum Clique
bitset<kN> a[kN], cs[kN];
                                                                     void add_edge( int vi , int ui , double ci )
{ e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
int ans, sol[kN], q, cur[kN], d[kN], n;
void init(int _n) {
 n = n, ans q = 0;
                                                                      for(int i=0; i<n; i++) d[0][i]=0;
                                                                      for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);</pre>
  for (int i = 0; i < n; i++) a[i].reset();</pre>
                                                                       for(int j=0; j<m; j++) {</pre>
void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
void csort(vector<int> &r, vector<int> &c) {
  int mx = 1, km = max(ans - q + 1, 1), t = 0,
                                                                        int v = e[j].v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
    m = int(r.size())
                                                                         d[i+1][u] = d[i][v]+e[j].c;
  cs[1].reset(); cs[2].reset();
for (int i = 0; i < m; i++) {
                                                                         prv[i+1][u] = v;
                                                                         prve[i+1][u] = j;
   int p = r[i], k = 1;
   while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
```

for (int i = 0; i < n; i++)

```
double solve(){
                                                                   ans = min(ans, dp[(1 << t) - 1][i]);
  // returns inf if no cycle, mmc otherwise
                                                                  return ans:
                                                                 }
  double mmc=inf;
  int st = -1;
                                                               } solver;
 bellman_ford();
                                                                3.14 Mo's Algorithm on Tree
  for(int i=0; i<n; i++) {</pre>
   double avg=-inf;
                                                               int q; vector< int > G[N];
   for(int k=0; k<n; k++) {</pre>
                                                                struct Que{
    if(d[n][i]<inf-eps)</pre>
                                                                 int u, v, id;
     avg=max(avg,(d[n][i]-d[k][i])/(n-k));
                                                                } que[ N ];
    else avg=max(avg,inf);
                                                               int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
void dfs( int u, int f ) {
   dfn[ u ] = dfn_++; int saved_rbp = stk_;
  if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
                                                                 for ( int v : G[ u ] ) {
 FZ(vst);edgeID.clear();cycle.clear();rho.clear();
                                                                  if ( v == f ) continue;
  for (int i=n; !vst[st]; st=prv[i--][st]) {
                                                                  dfs( v, u );
   vst[st]++;
                                                                  if ( stk_ - saved_rbp < SQRT_N ) continue;</pre>
   edgeID.PB(prve[i][st]);
                                                                 for ( ++ block_ ; stk_ != saved_rbp ; )
block_id[ stk[ -- stk_ ] ] = block_;
  rho.PB(st);
 while (vst[st] != 2) {
                                                                 stk[ stk_ ++ ] = u;
  int v = rho.back(); rho.pop_back();
   cycle.PB(v);
                                                               bool inPath[ N ];
  vst[v]++;
                                                                void Diff( int u ) {
                                                                if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
 reverse(ALL(edgeID));
                                                                else { /*add this edge*/ }
 edgeID.resize(SZ(cycle));
  return mmc;
                                                                void traverse( int& origin_u, int u ) {
}
                                                                for ( int g = lca( origin_u, u )
} mmc;
                                                                  origin_u != g ; origin_u = parent_of[ origin_u ] )
                                                                   Diff( origin_u );
     Minimum Steiner Tree
3.13
                                                                 for (int v = u; v != origin_u; v = parent_of[v])
// Minimum Steiner Tree
                                                                 Diff( v );
// 0(V 3^T + V^2 2^T)
                                                                 origin_u = u;
struct SteinerTree {
#define V 33
                                                                void solve() {
#define T 8
                                                                dfs( 1, 1 );
while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
#define INF 1023456789
int n, dst[V][V], dp[1 << T][V], tdst[V];</pre>
                                                                 sort( que, que + q, [](const Que& x, const Que& y) {
void init(int _n) {
                                                                  return tie( block_id[ x.u ], dfn[ x.v ] )
 n = _n;
                                                                      < tie( block_id[ y.u ], dfn[ y.v ] );
 for (int i = 0; i < n; i++) {</pre>
  for (int j = 0; j < n; j++)
dst[i][j] = INF * (i != j);</pre>
                                                                int'U = 1, V = 1;

for ( int i = 0 ; i < q ; ++ i ) {

pass( U, que[ i ].u );
 }
}
                                                                 pass( V, que[ i ].v );
void add_edge(int ui, int vi, int wi) {
                                                                  // we could get our answer of que[ i ].id
 dst[ui][vi] = min(dst[ui][vi], wi);
  dst[vi][ui] = min(dst[vi][ui], wi);
}
void shortest_path() {
                                                               Method 2:
 for (int k = 0; k < n; k++)
                                                                dfs u:
  for (int i = 0; i < n; i++)
                                                                 push u
    for (int j = 0; j < n; j++)
                                                                 iterate subtree
     dst[i][j] = min(dst[i][j], dst[i][k] + dst[k][j]);
                                                                 push u
                                                                Let P = LCA(u, v), and St(u) \le St(v)
int solve(const vector<int> &ter) {
                                                                if (P == u) query[St(u), St(v)]
 int t = (int)ter.size();
                                                               else query[Ed(u), St(v)], query[St(P), St(P)]
  for (int i = 1; i < (1 << t); i++)
                                                               */
   fill_n(dp[i], n, INF);
  fill_n(dp[0], n, 0);
                                                                3.15
                                                                      Tree Hashing
  for (int msk = 1; msk < (1 << t); msk++) {</pre>
                                                                uint64_t hsah(int u, int f) {
   if (msk == (msk & (-msk))) {
                                                                 uint64_t r = 127;
    int who = __lg(msk);
                                                                 for (int v : G[ u ]) if (v != f) {
    for (int i = 0; i < n; i++)
                                                                 uint64_t hh = hsah(v, u)
     dp[msk][i] = dst[ter[who]][i];
                                                                  r=(r+(hh*hh)%1010101333)%1011820613;
    continue;
                                                                 return r;
   for (int i = 0; i < n; i++)</pre>
    for (int submsk = (msk - 1) & msk; submsk; submsk =
     (submsk - 1) & msk)
                                                                3.16 Virtural Tree
     dp[msk][i] = min(dp[msk][i], dp[submsk][i] + dp[
    msk ^ submsk][i]);
                                                               inline bool cmp(const int &i, const int &j) {
   for (int i = 0; i < n; i++) {</pre>
                                                                 return dfn[i] < dfn[j];</pre>
    tdst[i] = INF
    for (int j = 0; j < n; j++)
                                                                void build(int vectrices[], int k) {
                                                                static int stk[MAX_N];
     tdst[i] = min(tdst[i], dp[msk][j] + dst[j][i]);
                                                                 sort(vectrices, vectrices + k, cmp);
  copy_n(tdst, n, dp[msk]);
                                                                 stk[sz++] = 0;
                                                                 for (int i = 0; i < k; ++i) {
  int u = vectrices[i], lca = LCA(u, stk[sz - 1]);</pre>
  int ans = INF;
                                                                  if (lca == stk[sz - 1]) stk[sz++] = u;
```

```
else {
  while (sz >= 2 && dep[stk[sz - 2]] >= dep[lca]) {
    addEdge(stk[sz - 2], stk[sz - 1]);
    sz--;
  }
  if (stk[sz - 1] != lca) {
    addEdge(lca, stk[--sz]);
    stk[sz++] = lca, vectrices[cnt++] = lca;
  }
  stk[sz++] = u;
}
  for (int i = 0; i < sz - 1; ++i)
  addEdge(stk[i], stk[i + 1]);
}</pre>
```

4 Matching & Flow

4.1 Bipartite Matching

```
struct BipartiteMatching {
vector<int> X[N];
 int fX[N], fY[N], n;
bitset<N> vis;
bool dfs(int x)
  for (auto i:X[x]) {
  if (vis[i]) continue;
   vis[i] = true;
   if (fY[i]==-1 || dfs(fY[i])){
    fY[fX[x] = i] = x;
    return true;
  }
  return false;
 void init(int n_, int m) {
 vis.reset();
 fill(X, X + (n = n_{-}), vector<int>());
 memset(fX, -1, sizeof(int) * n);
memset(fY, -1, sizeof(int) * m);
void add_edge(int x, int y){
 X[x].push_back(y); }
int solve() { // return how many pair matched
 int cnt = 0;
 for(int i=0;i<n;i++) {</pre>
  vis.reset()
  cnt += dfs(i);
 }
  return cnt;
```

4.2 Dijkstra Cost Flow

```
// kN = #(vertices)
// MCMF.{Init, AddEdge, MincostMaxflow}
// MincostMaxflow(source, sink, flow_limit, &cost)
// => flow
using Pii = pair<int, int>;
constexpr int kInf = 0x3f3f3f3f, kN = 500;
struct Edge {
int to, rev, cost, flow;
struct MCMF { // 0-based
int n{}, m{}, s{}, t{}
vector<Edge> graph[kN];
// Larger range for relabeling
int64_t dis[kN] = {}, h[kN] = {};
int p[kN] = {};
void Init(int nn) {
 n = nn;
 for (int i = 0; i < n; i++) graph[i].clear();</pre>
void AddEdge(int u, int v, int f, int c) {
 graph[u].push_back({v,
  static_cast<int>(graph[v].size()), c, f});
 graph[v].push_back(
   {u, static_cast<int>(graph[u].size()) - 1,
    -c, 0});
bool Dijkstra(int &max_flow, int64_t &cost) {
 priority_queue<Pii, vector<Pii>, greater<>> pq;
```

```
fill_n(dis, n, kInf);
  dis[s] = 0;
  pq.emplace(0, s);
  while (!pq.empty()) {
   auto u = pq.top();
   pq.pop();
   int v = u.second;
   if (dis[v] < u.first) continue;</pre>
   for (auto &e : graph[v]) {
    auto new_dis =
     dis[v] + e.cost + h[v] - h[e.to];
    if (e.flow > 0 && dis[e.to] > new_dis) {
     dis[e.to] = new_dis;
     p[e.to] = e.rev
     pq.emplace(dis[e.to], e.to);
  if (dis[t] == kInf) return false;
  for (int i = 0; i < n; i++) h[i] += dis[i];</pre>
  int d = max_flow;
  for (int u = t; u != s;
     u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   d = min(d, graph[e.to][e.rev].flow);
  max_flow -= d;
  cost += int64_t(d) * h[t];
  for (int u = t; u != s;
    u = graph[u][p[u]].to) {
   auto &e = graph[u][p[u]];
   e.flow += d:
   graph[e.to][e.rev].flow -= d;
  return true;
 int MincostMaxflow(
  int ss, int tt, int max_flow, int64_t &cost) {
  this->s = ss, this->t = tt;
  cost = 0;
  fill_n(h, n, 0);
  auto orig_max_flow = max_flow;
  while (Dijkstra(max_flow, cost) && max_flow) {}
  return orig_max_flow - max_flow;
};
4.3 Dinic
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
    int to, rev;
    Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS(){
    lv.assign(n, -1);
    queue<int> bfs;
    bfs.push(st); lv[st] = 0;
    while (not bfs.empty()){
      int u = bfs.front(); bfs.pop();
      for (auto e: G[u]) {
        if (e.cap <= 0 or lv[e.to]!=-1) continue;
        bfs.push(e.to); lv[e.to] = lv[u] + 1;
    return lv[ed] != -1;
  Cap DFS(int u, Cap f){
    if (u == ed) return f;
    Cap ret = 0:
    for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
```

auto &e = G[u][i];

if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>

Cap nf = DFS(e.to, min(f, e.cap));

ret += nf; e.cap -= nf; f -= nf;

G[e.to][e.rev].cap += nf;

if (f == 0) return ret;

```
if (ret == 0) lv[u] = -1;
     return ret;
public:
   void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
  G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
  st = st_, ed = ed_; Cap ret = 0;
     while (BFS()) {
        idx.assign(n, 0);
        Cap f = DFS(st, numeric_limits<Cap>::max());
        ret += f;
        if (f == 0) break;
      return ret;
   }
};
```

Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}.$
 - 2. For each edge (x,y,l,u), connect x o y with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum
 - of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t\to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f
 eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the
 - maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect t o s with capacity ∞ and let the flow from S to T be f'. If $f+f'
 eq \sum_{v\in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge \boldsymbol{e} on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipartite graph(X, Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M, x \to y$ otherwise. 2. DFS from unmatched vertices in X.

 - 3. $x \in X$ is chosen iff x is unvisited. 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect $x \to y$ with (cost,cap) = (c,1) if c > 0, otherwise connect $y \to x$ with (cost, cap) = (-c, 1)
 - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost, cap)=
 - (0, d(v))5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =(0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph

 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$ 2. Construct a max flow model, let ${\cal K}$ be the sum of all weights
 - Connect source $s \to v$, $v \in G$ with capacity K
 - 4. For each edge (u, v, w) in G, connect $u \to v$ and $v \to u$ with capacity
 - 5. For $v \in \mathit{G}$, connect it with sink $v \rightarrow t$ with capacity K + 2T - $\left(\sum_{e \in E(v)} w(e)\right) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.

 2. Create edge (u,v) with capacity w with w being the cost of choosing
 - u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of
- 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with ca-
- 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

```
4.5 General Graph Matching
```

```
namespace matching {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
 for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
 for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
 g[u].push_back(v);
 g[v].push_back(u);
int Find(int u) {
 return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
 static int tk = 0; tk++;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
   v[x] = tk;
   x = Find(pre[match[x]]);
void Blossom(int x, int y, int 1) {
 while (Find(x) != 1) {
 pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
if (fa[x] == x) fa[x] = 1;
  if (fa[y] == y) fa[y] = 1;
  x = pre[y];
bool Bfs(int r, int n) {
 for (int i = 0; i \le n; ++i) fa[i] = i, s[i] = -1;
 while (!q.empty()) q.pop();
 q.push(r);
 s[r] = 0;
 while (!q.empty()) {
  int x = q.front(); q.pop();
  for (int u : g[x]) {
   if (s[u] == -1)
    pre[u] = x, s[u] = 1;
    if (match[u] == n) {
     for (int a = u, b = x, last; b != n; a = last, b =
     pre[a])
      last = match[b], match[b] = a, match[a] = b;
     return true;
    q.push(match[u]);
    s[match[u]] = 0;
   } else if (!s[u] && Find(u) != Find(x)) {
    int 1 = LCA(u, x, n);
    Blossom(x, u, 1);
    Blossom(u, x, 1)
  }
 return false;
int Solve(int n) {
 int res = 0;
 for (int x = 0; x < n; ++x) {
  if (match[x] == n) res += Bfs(x, n);
 }
 return res;
     Global Min-Cut
```

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
w[x][y] += c; w[y][x] += c;
pair<int, int> phase(int n) {
 memset(v, false, sizeof(v));
 memset(g, 0, sizeof(g));
```

if (!v1[x] && d > s1k[x]) d = s1k[x];

```
int s = -1, t = -1;
                                                                   for (int x = 0; x < n; ++x) {
 while (true) {
                                                                    if (v1[x]) h1[x] += d;
                                                                    else slk[x] -= d;
  int c = -1;
  for (int i = 0; i < n; ++i) {</pre>
                                                                    if (vr[x]) hr[x] -= d;
   if (del[i] || v[i]) continue;
                                                                   for (int x = 0; x < n; ++x)
if (!v1[x] && !slk[x] && !check(x)) return;</pre>
   if (c == -1 \mid | g[i] > g[c]) c = i;
  if (c == -1) break;
 v[s = t, t = c] = true;
for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
                                                               public:
                                                                 void init( int n_ ) {
   g[i] += w[c][i];
                                                                  qu.resize(n = n_{-});
                                                                 fl.assign(n, -1); fr.assign(n, -1);
hr.assign(n, 0); hl.resize(n);
  }
                                                                 w.assign(n, vector<lld>(n));
 return make_pair(s, t);
                                                                 slk.resize(n); pre.resize(n);
int mincut(int n) {
                                                                 vl.resize(n); vr.resize(n);
 int cut = 1e9;
 memset(del, false, sizeof(del));
                                                                 void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 for (int i = 0; i < n - 1; ++i) {
                                                                 1ld solve() {
  int s, t; tie(s, t) = phase(n);
                                                                 for (int i = 0; i < n; ++i)
                                                                  hl[i] = *max_element(w[i].begin(), w[i].end());
  del[t] = true; cut = min(cut, g[t]);
  for (int j = 0; j < n; ++j) {
                                                                  for (int i = 0; i < n; ++i) bfs(i);</pre>
   w[s][j] += w[t][j]; w[j][s] += w[j][t];
                                                                 11d res = 0:
  }
                                                                  for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
                                                                  return res:
                                                                 }
 return cut;
                                                               } km;
                                                               4.9 Minimum Cost Circulation
4.7
     GomoryHu Tree
                                                               struct Edge { int to, cap, rev, cost; };
int g[maxn];
                                                               vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
vector<edge> GomoryHu(int n){
 vector<edge> rt;
 for(int i=1;i<=n;++i)g[i]=1;</pre>
                                                               bool mark[kN];
 for(int i=2;i<=n;++i){</pre>
                                                               int NegativeCycle(int n) {
                                                                memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
  int t=g[i];
  flow.reset(); // clear flows on all edge
  rt.push_back({i,t,flow(i,t)});
                                                                 int upd = -1;
 flow.walk(i); // bfs points that connected to i (use
  edges not fully flow)
                                                                 for (int i = 0; i <= n; ++i)
                                                                 for (int j = 0; j < n; ++j) {
  for(int j=i+1;j<=n;++j){</pre>
                                                                   int idx = 0;
                                                                   for (auto &e : g[j]) {
   if(g[j]==t && flow.connect(j))g[j]=i; // check if i
                                                                    if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
    can reach j
 }
                                                                     dist[e.to] = dist[j] + e.cost;
                                                                     pv[e.to] = j, ed[e.to] = idx;
 return rt;
                                                                     if (i == n) {
                                                                      upd = j;
                                                                      while(!mark[upd])mark[upd]=1,upd=pv[upd];
4.8 Kuhn Munkres
                                                                      return upd;
class KM {
                                                                     }
private:
 static constexpr 1ld INF = 1LL << 60;</pre>
                                                                    idx++;
 vector<lld> hl,hr,slk;
 vector<int> fl,fr,pre,qu;
                                                                  }
 vector<vector<lld>> w;
 vector<bool> v1,vr;
                                                                 return -1;
 int n, ql, qr;
 bool check(int x) {
                                                               int Solve(int n) {
  if (v1[x] = true, f1[x] != -1)
                                                                 int rt = -1, ans = 0;
   return vr[qu[qr++] = f1[x]] = true;
                                                                 while ((rt = NegativeCycle(n)) >= 0) {
  while (x != -1) swap(x, fr[fl[x] = pre[x]]);
                                                                 memset(mark, false, sizeof(mark));
                                                                  vector<pair<int, int>> cyc;
  return false;
                                                                  while (!mark[rt]) {
 void bfs(int s) {
                                                                   cyc.emplace_back(pv[rt], ed[rt]);
  fill(slk.begin(), slk.end(), INF);
                                                                   mark[rt] = true;
  fill(v1.begin(), v1.end(), false);
                                                                   rt = pv[rt];
  fill(vr.begin(), vr.end(), false);
  ql = qr = 0;
                                                                  reverse(cyc.begin(), cyc.end());
  vr[qu[qr++] = s] = true;
                                                                  int cap = kInf;
  while (true) {
                                                                  for (auto &i : cyc)
                                                                   auto &e = g[i.first][i.second];
   11d d;
   while (ql < qr) {</pre>
                                                                   cap = min(cap, e.cap);
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
     if(!v1[x]\&s1k[x]>=(d=h1[x]+hr[y]-w[x][y])){
                                                                  for (auto &i : cyc) {
                                                                   auto &e = g[i.first][i.second];
      if (pre[x] = y, d) slk[x] = d;
      else if (!check(x)) return;
                                                                   e.cap -= cap;
     }
                                                                   g[e.to][e.rev].cap += cap;
    }
                                                                   ans += e.cost * cap;
   d = INF;
   for (int x = 0; x < n; ++x)
                                                                 return ans;
```

4.10 Minimum Cost Maximum Flow

```
class MiniCostMaxiFlow{
 using Cap = int; using Wei = int64_t;
 using PCW = pair<Cap,Wei>;
 static constexpr Cap INF_CAP = 1 << 30;</pre>
 static constexpr Wei INF_WEI = 1LL<<60;</pre>
private:
 struct Edge{
  int to, back;
Cap cap; Wei wei;
  Edge() {}
 Edge(int a,int b, Cap c, Wei d):
   to(a),back(b),cap(c),wei(d) {}
 };
 int ori, edd;
 vector<vector<Edge>> G;
 vector<int> fa, wh;
 vector<bool> inq;
 vector<Wei> dis;
 PCW SPFA(){
  fill(inq.begin(),inq.end(),false)
  fill(dis.begin(), dis.end(), INF_WEI);
  queue<int> qq; qq.push(ori);
  dis[ori] = 0;
  while(not qq.empty()){
   int u=qq.front();qq.pop();
   inq[u] = false;
   for(int i=0;i<SZ(G[u]);++i){</pre>
    Edge e=G[u][i];
    int v=e.to; Wei d=e.wei;
    if(e.cap <= 0 | |dis[v] <= dis[u] + d)
     continue
    dis[v] = dis[u] + d;
    fa[v] = u, wh[v] = i;
if (inq[v]) continue;
    qq.push(v);
    inq[v] = true;
  if(dis[edd]==INF_WEI) return {-1, -1};
  Cap mw=INF_CAP;
  for(int i=edd;i!=ori;i=fa[i])
   mw=min(mw,G[fa[i]][wh[i]].cap);
  for (int i=edd;i!=ori;i=fa[i]){
   auto &eg=G[fa[i]][wh[i]];
   eg.cap -= mw;
   G[eg.to][eg.back].cap+=mw;
  return {mw, dis[edd]};
public:
 void init(int n){
  G.clear();G.resize(n);
  fa.resize(n);wh.resize(n);
  inq.resize(n); dis.resize(n);
 void add_edge(int st, int ed, Cap c, Wei w){
  G[st].emplace_back(ed,SZ(G[ed]),c,w);
  G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
 PCW solve(int a, int b){
 ori = a, edd = b;
  Cap cc=0; Wei ww=0;
  while(true)
  PCW ret=SPFA();
   if(ret.first==-1) break;
   cc+=ret.first;
   ww+=ret.first * ret.second;
  }
  return {cc,ww};
 }
} mcmf;
```

4.11 Maximum Weight Graph Matching

```
struct WeightGraph {
  static const int inf = INT_MAX;
  static const int maxn = 514;
  struct edge {
  int u, v, w;
  edge(){}
  edge(int u, int v, int w): u(u), v(v), w(w) {}
```

```
int n, n_x;
edge g[maxn * 2][maxn * 2];
int lab[maxn * 2];
int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa
   [maxn * 2];
int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
   maxn * 21
vector<int> flo[maxn * 2];
queue<int> q;
int e_delta(const edge &e) { return lab[e.u] + lab[e.v
   ] - g[e.u][e.v].w * 2; }
void update_slack(int u, int x) { if (!slack[x] ||
   e_delta(g[u][x]) < e_delta(g[slack[x]][x])) slack[x]
   ] = u; }
void set_slack(int x) {
 slack[x] = 0;
 for (int u = 1; u <= n; ++u)
  if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
   update_slack(u, x);
void q_push(int x) {
if (x \le n) q.push(x);
 else for (size_t i = 0; i < flo[x].size(); i++)</pre>
   q_push(flo[x][i]);
void set_st(int x, int b) {
st[x] = b;
 if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)</pre>
    set_st(flo[x][i], b);
int get_pr(int b, int xr) {
 int pr = find(flo[b].begin(), flo[b].end(), xr) - flo
   [b].begin();
 if (pr % 2 == 1) {
  reverse(flo[b].begin() + 1, flo[b].end());
 return (int)flo[b].size() - pr;
}
 return pr;
void set_match(int u, int v) {
match[u] = g[u][v].v;
 if (u <= n) return;</pre>
 edge e = g[u][v];
 int xr = flo_from[u][e.u], pr = get_pr(u, xr)
 for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo</pre>
   [u][i ^ 1]);
set_match(xr, v);
rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
   end());
void augment(int u, int v) {
for (; ; ) {
  int xnv = st[match[u]];
  set_match(u, v);
 if (!xnv) return
  set_match(xnv, st[pa[xnv]]);
  u = st[pa[xnv]], v = xnv;
int get_lca(int u, int v) {
 static int t = 0;
 for (++t; u || v; swap(u, v)) {
 if (u == 0) continue;
if (vis[u] == t) return u;
 vis[u] = t;
 u = st[match[u]];
 if (u) u = st[pa[u]];
 return 0;
void add_blossom(int u, int lca, int v) {
 int b = n + 1;
 while (b <= n_x && st[b]) ++b;</pre>
 if (b > n_x) ++n_x;
 lab[b] = 0, S[b] = 0;
 match[b] = match[lca];
 flo[b].clear();
 flo[b].push_back(lca);
 for (int x = u, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
```

]]), q_push(y);

```
reverse(flo[b].begin() + 1, flo[b].end())
 for (int x = v, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[match[x
   ]]), q_push(y);
 set_st(b, b);
 for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].w
   = 0;
 for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
 for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  for (int x = 1; x <= n_x; ++x)
   if (g[b][x].w == 0 \mid \mid e_delta(g[xs][x]) < e_delta(g[xs][x])
   [b][x]))
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
  for (int x = 1; x <= n; ++x)
   if (flo_from[xs][x]) flo_from[b][x] = xs;
 set_slack(b);
}
void expand_blossom(int b) {
 for (size_t i = 0; i < flo[b].size(); ++i)
set_st(flo[b][i], flo[b][i]);</pre>
 int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b,
   xr);
 for (int i = 0; i < pr; i += 2) {
  int xs = flo[b][i], xns = flo[b][i + 1];</pre>
  pa[xs] = g[xns][xs].u;
  S[xs] = 1, S[xns] = 0;
  slack[xs] = 0, set_slack(xns);
  q_push(xns);
 S[xr] = 1, pa[xr] = pa[b];
 for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
  int xs = flo[b][i];
  S[xs] = -1, set_slack(xs);
 st[b] = 0;
bool on_found_edge(const edge &e) {
 int u = st[e.u], v = st[e.v];
 if (S[v] == -1) {
  pa[v] = e.u, S[v] = 1;
  int nu = st[match[v]];
  slack[v] = slack[nu] = 0;
 S[nu] = 0, q_push(nu);
} else if (S[v] == 0) {
  int lca = get_lca(u, v);
  if (!lca) return augment(u,v), augment(v,u), true;
  else add_blossom(u, lca, v);
 return false;
bool matching() {
 memset(S + 1, -1, sizeof(int) * n_x;
memset(slack + 1, 0, sizeof(int) * n_x);
 q = queue<int>();
 for (int x = 1; x <= n_x; ++x)
 if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0,
   q_push(x);
 if (q.empty()) return false;
 for (; ; ) {
  while (q.size()) {
   int u = q.front(); q.pop();
   if (S[st[u]] == 1) continue;
   for (int v = 1; v <= n; ++v)
    if (g[u][v].w > 0 && st[u] != st[v]) {
     if (e_delta(g[u][v]) == 0) {
      if (on_found_edge(g[u][v])) return true;
     } else update_slack(u, st[v]);
    }
  int d = inf;
  for (int b = n + 1; b \le n_x; ++b)
   if (st[b] == b && S[b] == 1) d = min(d, lab[b] / 2)
  for (int x = 1; x <= n_x; ++x)
   if (st[x] == x && slack[x]) {
  if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x])
    else if (S[x] == 0) d = min(d, e_delta(g[slack[x
   ]][x]) / 2);
```

```
for (int u = 1; u <= n; ++u) {
    if (S[st[u]] == 0) {
     if (lab[u] <= d) return 0;</pre>
      lab[u] -= d;
    } else if (S[st[u]] == 1) lab[u] += d;
   for (int b = n + 1; b \le n_x; ++b)
    if (st[b] == b) {
      if (S[st[b]] == 0) lab[b] += d * 2;
     else if (S[st[b]] == 1) lab[b] -= d * 2;
    q = queue<int>();
   for (int x = 1; x <= n_x; ++x)
if (st[x] == x && slack[x] && st[slack[x]] != x &&
     e_delta(g[slack[x]][x]) == 0)
     if (on_found_edge(g[slack[x]][x])) return true;
    for (int b = n + 1; b <= n_x; ++b)
    if (st[b] == b && S[b] == 1 && lab[b] == 0)
     expand_blossom(b);
  return false;
 pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n_x = n:
  int n_matches = 0;
  long long tot_weight = 0;
  for (int u = 0; u \le n; t+u) st[u] = u, flo[u].clear
     ();
  int w_max = 0;
  for (int u = 1; u <= n; ++u)
   for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);</pre>
    w_{max} = max(w_{max}, g[u][v].w);
  for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
  while (matching()) ++n_matches;
  for (int u = 1; u <= n; ++u)
   if (match[u] && match[u] < u)</pre>
    tot_weight += g[u][match[u]].w;
  return make_pair(tot_weight, n_matches);
 void add_edge(int ui, int vi, int wi) { g[ui][vi].w =
     g[vi][ui].w = wi; }
 void init(int _n) {
  n = _n;
  for (int u = 1; u <= n; ++u)
   for (int v = 1; v <= n; ++v)
     g[u][v] = edge(u, v, 0);
};
```

5 Math

5.1 $\lfloor \frac{n}{i} \rfloor$ Enumeration

 $T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor$

5.2 Strling Number

5.2.1 First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1) \dots (x+n-1) = \sum_{k=0}^n S_1(n,k) x^k$$

$$g(x) = x(x+1) \dots (x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.2.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

```
S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}
```

5.3 ax+by=gcd

```
// ax+ny = 1, ax+ny == ax == 1 \pmod{n}
void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
if (y == 0) g=x, a=1, b=0;
else exgcd(y, x\%y, g, b, a), b=(x/y)*a;
```

5.4 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
vector<T> d(output.size() + 1), me, he;
for (size_t f = 0, i = 1; i <= output.size(); ++i) {</pre>
  for (size_t j = 0; j < me.size(); ++j)</pre>
  d[i] += output[i - j - 2] * me[j];
if ((d[i] -= output[i - 1]) == 0) continue;
  if (me.empty()) {
   me.resize(f = i);
   continue;
 vector<T> o(i - f - 1);
T k = -d[i] / d[f]; o.push_back(-k);
 for (T x : he) o.push_back(x * k);
 if (o.size() < me.size()) o.resize(me.size());</pre>
  for (size_t j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
  if (i-f+he.size() >= me.size()) he = me, f = i;
 me = o:
return me;
```

5.5 Charateristic Polynomial

```
vector<vector<int>> Hessenberg(const vector<vector<int</pre>
    >> &A) {
 int N = A.size();
vector<vector<int>> H = A;
for (int i = 0; i < N - 2; ++i) {
 if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {
    if (H[j][i]) {
     for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
    ][k]);
     for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k
    ][j]);
    }
   }
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
 for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
   * H[i + 1][k] * (kP - coef)) % kP;
for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
return H:
vector<int> CharacteristicPoly(const vector<vector<int
    >> &A) {
int N = A.size();
 auto H = Hessenberg(A);
for (int i = 0; i < N; ++i) {
 for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
vector<vector<int>> P(N + 1, vector<int>(N + 1));
P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
 P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
    1];
  int val = 1;
 for (int j = i - 1; j >= 0; --j) {
  int coef = 1LL * val * H[j][i - 1] % kP;
   for (int k = 0; k \le j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
```

```
if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
 if (N & 1) {
  for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];</pre>
 return P[N];
5.6 Chinese Remainder
1ld crt(lld ans[], lld pri[], int n){
 lld M = 1, ret = 0;
 for(int i=0;i<n;i++) M *= pri[i];</pre>
 for(int i=0;i<n;i++){</pre>
  lld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
  ret += (ans[i]*(M/pri[i])%M * iv)%M;
  ret %= M;
 return ret;
}
/*
Another:
x = a1 \% m1
x = a2 \% m2
g = gcd(m1, m2)
assert((a1-a2)%g==0)
[p, q] = exgcd(m2/g, m1/g)
return a2+m2*(p*(a1-a2)/g)
0 <= x < lcm(m1, m2)
5.7 De-Bruijn
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
 if (t > n) {
  if (n % p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
 } else {
  aux[t] = aux[t - p];
  db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {
   aux[t] = i;
   db(t + 1, t, n, k);
  }
int de_bruijn(int k, int n) {
 // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
 if (k == 1) {
  res[0] = 0;
  return 1:
 for (int i = 0; i < k * n; i++) aux[i] = 0;
 sz = 0;
 db(1, 1, n, k);
 return sz;
5.8 DiscreteLog
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
  // x^? \equiv y (mod M)
  Int t = 1, c = 0, g = 1;
  for (Int M_ = M; M_ > 0; M_ >>= 1)
    g = g * x % M;
  for (g = gcd(g, M); t % g != 0; ++c) {
    if (t == y) return c;
    t = t * x % M;
  if (y % g != 0) return -1;
  t /= g, y /= g, M /= g;
  Int h = 0, gs = 1;
for (; h * h < M; ++h) gs = gs * x % M;
  unordered_map<Int, Int> bs;
  for (Int s = 0; s < h; bs[y] = ++s)
    y = y * x % M;
  for (Int s = 0; s < M; s += h) {
```

t = t * gs % M;

if (bs.count(t)) return c + s + h - bs[t];

5.10 ExtendedFloorSum

```
\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \\ - \frac{\frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1)}{-h(c, c-b-1, a, m-1)),} & \text{otherwise} \end{cases} \\ h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{c} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}
```

5.11 Fast Fourier Transform

```
const int mod = 1000000007:
const int M1 = 985661441; // G = 3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
  static_assert (M1 <= M2 && M2 <= M3);
 constexpr int64_t r12 = modpow(M1, M2-2, M2);
 constexpr int64_t r13 = modpow(M1, M3-2, M3);
 constexpr int64_t r23 = modpow(M2, M3-2, M3);
 constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
 B = (B - A + M2) * r12 % M2;
 C = (C - A + M3) * r13 % M3;
 C = (C - B + M3) * r23 % M3;
  return (A + B * M1 + C * M1M2) % mod;
namespace fft {
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
for (int i = 0; i <= maxn; i++)</pre>
 omega[i] = cplx(cos(2 * pi * j / maxn),
     sin(2 * pi * j / maxn));
void fft(vector<cplx> &v, int n) {
int z = __builtin_ctz(n) - 1;
for (int i = 0; i < n; ++i) {
 int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^{=(i >> j & 1)<<(z - j);
 if (x > i) swap(v[x], v[i]);
for (int s = 2; s <= n; s <<= 1) {
 int z = s >> 1;
for (int i = 0; i < n; i += s) {</pre>
  for (int k = 0; k < z; ++k) {
  cplx x = v[i + z + k] * omega[maxn / s * k];
    v[i + z + k] = v[i + k] - x;
    v[i + k] = v[i + k] + x;
void ifft(vector<cplx> &v, int n) {
fft(v, n); reverse(v.begin() + 1, v.end());
for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);
VL convolution(const VI &a, const VI &b) {
// Should be able to handle N <= 10^5, C <= 10^4
```

```
int sz = 1;
 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
 vector<cplx> v(sz);
 for (int i = 0; i < sz; ++i) {
  double re = i < a.size() ? a[i] : 0;</pre>
  double im = i < b.size() ? b[i] : 0;</pre>
  v[i] = cplx(re, im);
 fft(v, sz);
 for (int i = 0; i <= sz / 2; ++i) {
int j = (sz - i) & (sz - 1);
  cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
  * cplx(0, -0.25);
if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
     ].conj()) * cplx(0, -0.25);
  v[i] = x;
 ifft(v, sz);
 VL c(sz);
 for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
 return c:
VI convolution_mod(const VI &a, const VI &b, int p) {
 int sz = 1:
 while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
 vector<cplx> fa(sz), fb(sz);
for (int i = 0; i < (int)a.size(); ++i)</pre>
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
 for (int i = 0; i < (int)b.size(); ++i)
fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
 fft(fa, sz), fft(fb, sz);
 double r = 0.25 / sz;
 cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
 for (int i = 0; i <= (sz >> 1); ++i) {
  int j = (sz - i) & (sz - 1);
  cplx a1 = (fa[i] + fa[j].conj());
  cplx a2 = (fa[i] - fa[j].conj()) * r2;
  cplx b1 = (fb[i] + fb[j].conj()) * r3;
  cplx b2 = (fb[i] - fb[j].conj()) * r4;
  if (i != j) {
   cplx c1 = (fa[j] + fa[i].conj());

cplx c2 = (fa[j] - fa[i].conj()) * r2;

cplx d1 = (fb[j] + fb[i].conj()) * r3;
   cplx d2 = (fb[j] - fb[i].conj()) * r4;
   fa[i] = c1 * d1 + c2 * d2 * r5;
fb[i] = c1 * d2 + c2 * d1;
  fa[j] = a1 * b1 + a2 * b2 * r5;
  fb[j] = a1 * b2 + a2 * b1;
 fft(fa, sz), fft(fb, sz);
 vector<int> res(sz);
 for (int i = 0; i < sz; ++i) {
  long long a = round(fa[i].re), b = round(fb[i].re),
        c = round(fa[i].im);
  res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
 }
 return res;
}}
5.12 FloorSum
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
// @return sum_\{i=0\}^{n-1} floor((ai + b)/m) mod 2^64
1lu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
 llu ans = 0;
 while (true) {
  if (a >= m) {
   ans += n * (n - 1) / 2 * (a / m); a %= m;
  if (b >= m) {
   ans += n * (b / m); b %= m;
  llu y_max = a * n + b;
  if (y_max < m) break;</pre>
  // y_{max} < m * (n + 1)
  // floor(y_max / m) <= n
  n = (11u)(y_max / m), b = (11u)(y_max % m);
  swap(m, a);
 return ans;
```

```
11d floor_sum(11d n, 11d m, 11d a, 11d b) {
                                                                  if(x<2)return 0;</pre>
llu ans = 0;
                                                                  if(!(x&1))return x==2;
if (a < 0) {
                                                                  llu x1=x-1; int t=0;
 11u \ a2 = (a \% m + m) \% m;
                                                                  while(!(x1&1))x1>>=1,t++;
 ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
                                                                  for(llu m:magic)if(witn(m,x1,x,t))return 0;
if (b < 0) {
                                                                 5.16 NTT
 11u b2 = (b \% m + m) \% m;
 ans -= 1ULL * n * ((b2 - b) / m);
                                                                 template <int mod, int G, int maxn>
 b = b2:
                                                                 struct NTT {
                                                                  static_assert (maxn == (maxn & -maxn));
return ans + floor_sum_unsigned(n, m, a, b);
                                                                  int roots[maxn];
                                                                  NTT () {
                                                                   int r = modpow(G, (mod - 1) / maxn);
5.13 FWT
                                                                   for (int i = maxn >> 1; i; i >>= 1) {
                                                                    roots[i] = 1;
/* xor convolution:
                                                                    for (int j = 1; j < i; j++)
roots[i + j] = modmul(roots[i + j - 1], r);</pre>
* x = (x0,x1) , y = (y0,y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                    r = modmul(r, r);
* x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
* z = (1/2) * z''
                                                                   }
                                                                  // n must be 2^k, and 0 \le F[i] < mod
                                                                  void operator()(int F[], int n, bool inv = false) {
  for (int i = 0, j = 0; i < n; i++) {</pre>
* or convolution:
* x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                                    if (i < j) swap(F[i], F[j]);</pre>
 * and convolution:
                                                                    for (int k = n>1; (j^=k) < k; k>>=1);
* x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
const LL MOD = 1e9+7;
                                                                   for (int s = 1; s < n; s *= 2) {
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
                                                                    for (int i = 0; i < n; i += s * 2) {
for( int d = 1 ; d < N ; d <<= 1 ) {</pre>
                                                                     for (int j = 0; j < s; j++) {
  int a = F[i+j];</pre>
  int d2 = d << 1;
  for( int s = 0 ; s < N ; s += d2 )
   for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
  LL ta = x[ i ] , tb = x[ j ];
  x[ i ] = ta+tb;</pre>
                                                                      int b = modmul(F[i+j+s], roots[s+j]);
                                                                      F[i+j] = modadd(a, b); // a + b
                                                                      F[i+j+s] = modsub(a, b); // a - b
    x[ j ] = ta-tb;
    if( x[ i ] >= MOD ) x[ i ] -= MOD;
    if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
                                                                   if (inv) {
   }
                                                                    int invn = modinv(n);
                                                                    for (int i = 0; i < n; i++)</pre>
if( inv )
                                                                     F[i] = modmul(F[i], invn);
 for( int i = 0 ; i < N ; i++ ) {</pre>
                                                                    reverse(F + 1, F + n);
   x[ i ] *= inv( N, MOD );
   x[ i ] %= MOD;
                                                                  }
  }
                                                                 NTT<2013265921, 31, 1048576> ntt;
5.14
      Gauss Elimination
                                                                 5.17 Range Sieve
void gauss(vector<vector<double>> &d) {
                                                                 const int MAX_SQRT_B = 50000;
int n = d.size(), m = d[0].size();
                                                                 const int MAX_L = 200000 + 5;
for (int i = 0; i < m; ++i) {
                                                                 bool is_prime_small[MAX_SQRT_B], is_prime[MAX_L];
 int p = -1;
                                                                 void sieve(lld 1, lld r){ // [1, r)
  for (int j = i; j < n; ++j) {
                                                                  for(lld i=2;i*i<r;i++) is_prime_small[i] = true;</pre>
   if (fabs(d[j][i]) < eps) continue;</pre>
   if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
                                                                  for(lld i=1;i<r;i++) is_prime[i-1] = true;</pre>
                                                                  if(l==1) is_prime[0] = false;
                                                                  for(1ld i=2;i*i<r;i++){</pre>
 if (p == -1) continue;
 for (int j = 0; j < m'; ++j) swap(d[p][j], d[i][j]); for (int j = 0; j < n; ++j) {
                                                                   if(!is_prime_small[i]) continue;
                                                                   for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;</pre>
                                                                   for(lld j=std::max(2LL, (l+i-1)/i)*i;j<r;j+=i)</pre>
  if (i == j) continue;
                                                                    is_prime[j-1]=false;
   double z = d[j][i] / d[i][i];
   for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
                                                                 5.18 Partition Number
                                                                 int b = sqrt(n);
5.15
      Miller Rabin
                                                                 ans[0] = tmp[0] = 1;
                                                                 for (int i = 1; i <= b; i++) {
bool isprime(llu x){
                                                                  for (int rep = 0; rep < 2; rep++)</pre>
static llu magic[]={2,325,9375,28178,\
                                                                   for (int j = i; j <= n - i * i; j++)
          450775,9780504,1795265022};
                                                                    modadd(tmp[j], tmp[j-i]);
static auto witn=[](llu a,llu u,llu n,int t)
                                                                  for (int j = i * i; j <= n; j++)
modadd(ans[j], tmp[j - i * i]);</pre>
 ->bool{
 if (!(a = mpow(a%n,u,n)))return 0;
  while(t--){
  llu a2=mul(a,a,n);
                                                                 5.19 Pi Count (Linear Sieve)
   if(a2==1 && a!=1 && a!=n-1)
                                                                 static constexpr int N = 1000000 + 5;
    return 1;
   a = a2;
                                                                 11d pi[N];
  }
                                                                 vector<int> primes;
                                                                 bool sieved[N];
  return a!=1;
};
                                                                 11d cube_root(11d x){
```

Poly &isz(int _n) { return resize(_n), *this; }

```
lld s=cbrt(x-static_cast<long double>(0.1));
                                                                Poly &iadd(const Poly &rhs) { // n() == rhs.n()
                                                                fi(0, size())(*this)[i] = modadd((*this)[i], rhs[i]);
while(s*s*s <= x) ++s;</pre>
                                                                 return *this;
return s-1;
11d square_root(11d x){
                                                                Poly &imul(int k) {
                                                                 fi(0, size())(*this)[i] = modmul((*this)[i], k);
1ld s=sqrt(x-static_cast<long double>(0.1));
while(s*s <= x) ++s;
                                                                 return *this;
return s-1;
                                                                Poly Mul(const Poly &rhs) const {
                                                                const int _n = n2k(size() + rhs.size() - 1);
Poly X(*this, _n), Y(rhs, _n);
void init(){
primes.reserve(N);
                                                                 ntt(X.data(), _n), ntt(Y.data(), _n)
fi(0, _n) X[i] = modmul(X[i], Y[i]);
primes.push_back(1);
for(int i=2;i<N;i++) {</pre>
 if(!sieved[i]) primes.push_back(i);
                                                                 ntt(X.data(), _n, true)
 pi[i] = !sieved[i] + pi[i-1];
                                                                 return X.isz(size() + rhs.size() - 1);
 for(int p: primes) if(p > 1) {
  if(p * i >= N) break;
                                                                Poly Inv() const { // coef[0] != 0
                                                                 if (size() == 1) return V{modinv(*begin())};
  sieved[p * i] = true;
  if(p % i == 0) break;
                                                                 const int _n = n2k(size() * 2);
                                                                 Poly Xi = Poly(*this, (size() + 1) / 2).Inv().isz(_n)
                                                                     Y(*this, _n);
                                                                ntt(Xi.data(), _n), ntt(Y.data(), _n);
fi(0, _n) Xi[i] = modmul(Xi[i], modsub(2, modmul(Xi[i], Y[i])));
11d phi(11d m, 11d n) {
static constexpr int MM = 80000, NN = 500;
static 1ld val[MM][NN];
                                                                 ntt(Xi.data(), _n, true);
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
                                                                 return Xi.isz(size());
if(n == 0) return m;
if(primes[n] >= m) return 1;
                                                                Poly Sqrt() const { // coef[0] \in [1, mod)^2
1ld ret = phi(m,n-1)-phi(m/primes[n],n-1);
                                                                 if (size() == 1) return V{QuadraticResidue((*this)
                                                                   [0], mod)};
if(m<MM&&n<NN) val[m][n] = ret+1;
                                                                 Poly X = Poly(*this, (size() + 1) / 2).Sqrt().isz(
return ret:
                                                                   size());
1ld pi_count(1ld);
                                                                 return X.iadd(Mul(X.Inv()).isz(size())).imul(mod / 2
11d P2(11d m, 11d n) {
                                                                   + 1);
11d sm = square_root(m), ret = 0;
for(lld i = n+1;primes[i]<=sm;i++)</pre>
                                                                pair<Poly, Poly> DivMod(const Poly &rhs) const {
 ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
                                                                 if (size() < rhs.size()) return {V{0}, *this};</pre>
return ret;
                                                                 const int _n = size() - rhs.size() + 1;
                                                                 Poly X(rhs); X.irev().isz(_n);
                                                                 Poly Y(*this); Y.irev().isz(_n);
11d pi_count(11d m) {
if(m < N) return pi[m];</pre>
                                                                 Poly Q = Y.Mul(X.Inv()).isz(_n).irev();
                                                                 X = rhs.Mul(Q), Y = *this;
fi(0, size()) Y[i] = modsub(Y[i], X[i]);
11d n = pi_count(cube_root(m));
return phi(m, n) + n - 1 - P2(m, n);
                                                                 return {Q, Y.isz(max<int>(1, rhs.size() - 1))};
5.20 Pollard Rho
                                                                Poly Dx() const {
// does not work when n is prime
                                                                 Poly ret(size() - 1);
                                                                 fi(0, ret.size()) ret[i] = modmul(i + 1, (*this)[i +
// return any non-trivial factor
llu pollard_rho(llu n){
static auto f=[](llu x,llu k,llu m){
                                                                 return ret.isz(max<int>(1, ret.size()));
  return add(k,mul(x,x,m),m);};
if (!(n&1)) return 2;
mt19937 rnd(120821011);
                                                                Poly Sx() const {
                                                                Poly ret(size() + 1);
while(true){
                                                                 fi(0, size()) ret[i + 1] = modmul(modinv(i + 1), (*
 llu y=2,yy=y,x=rnd()%n,t=1;
for(llu sz=2;t==1;sz<<=1,y=yy) {</pre>
                                                                   this)[i]);
                                                                 return ret:
   for(llu i=0;t==1&&i<sz;++i){</pre>
   yy=f(yy,x,n);
                                                                Poly _tmul(int nn, const Poly &rhs) const {
                                                                Poly Y = Mul(rhs).isz(size() + nn - 1);
    t=gcd(yy>y?yy-y:y-yy,n);
                                                                 return V(Y.data() + size() - 1, Y.data() + Y.size());
                                                                V _eval(const V &x, const vector<Poly> &up) const {
  if(t!=1&&t!=n) return t;
                                                                 const int _n = (int)x.size();
                                                                 if (!_n) return {};
                                                                 vector<Poly> down(_n * 2);
5.21 Polynomial Operations
                                                                 // down[1] = DivMod(up[1]).second;
                                                                 // fi(2, _n * 2) down[i] = down[i / 2].DivMod(up[i]).
using V = vector<int>;
#define fi(s, n) for (int i = int(s); i < int(n); ++i)</pre>
                                                                   second;
template <int mod, int G, int maxn> struct Poly : V {
                                                                 down[1] = Poly(up[1]).irev().isz(size()).Inv().irev()
                                                                 ._tmul(_n, *this);
fi(2, _n * 2) down[i] = up[i ^ 1]._tmul(up[i].size()
static uint32_t n2k(uint32_t n) {
 if (n <= 1) return 1;
  return 1u << (32 - __builtin_clz(n - 1));</pre>
                                                                   - 1, down[i / 2]);
                                                                 V y(_n); fi(0, _n) y[i] = down[_n + i][0];
static NTT<mod,G,maxn> ntt; // coefficients in [0, P)
                                                                 return y;
 explicit Poly(int n = 1) : V(n) {}
Poly(const V &v) : V(v) {}
                                                                static vector<Poly> _tree1(const V &x) {
Poly(const Poly &p, size_t n) : V(n) {
                                                                 const int _n = (int)x.size();
                                                                 vector<Poly> up(_n * 2);
 copy_n(p.data(), min(p.size(), n), data());
                                                                 fi(0, _n) up[_n + i] = V\{modsub(mod, x[i]), 1\};
                                                                 for(int i=_n-1;i>0;--i) up[i] = up[i * 2].Mul(up[i *
Poly &irev() { return reverse(data(), data() + size())
      *this; }
                                                                   2 + 1]);
```

return up;

```
int w_{-} = w;
                                                                 if (w_{-} == -1) w_{-} = rhs.w;
V Eval(const V &x) const { return _eval(x, _tree1(x));
                                                                 assert(w_ != -1 and w_ == rhs.w);
                                                                 return { MOD, w_,
(x * rhs.x + y * rhs.y % MOD * w) % MOD,
 static Poly Interpolate(const V &x, const V &y) {
 const int _n = (int)x.size();
                                                                   (x * rhs.y + y * rhs.x) % MOD };
  vector<Poly> up = _{tree1(x), down(_n * 2);}
  V z = up[1].Dx()._eval(x, up);
 fi(0, _n) z[i] = modmul(y[i], modinv(z[i]));
 fi(0, _n) down[_n + i] = V{z[i]};
for(int i=_n-1;i>0;--i) down[i] = down[i * 2].Mul(up[
                                                               int get_root(int n, int P) {
                                                                 if (P == 2 or n == 0) return n;
                                                                 if (qpow(n, (P - 1) / 2, P) != 1) return -1;
    i * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
                                                                 auto check = [&](int x) {
  return down[1];
                                                                 return qpow(x, (P - 1) / 2, P); };
if (check(n) == P-1) return -1;
Poly Ln() const \{ // coef[0] == 1 \}
 return Dx().Mul(Inv()).Sx().isz(size());
                                                                 int64_t a; int w; mt19937 rnd(7122);
                                                                 do { a = rnd() % P;
  w = ((a * a - n) % P + P) % P;
Poly Exp() const \{ // coef[0] == 0 \}
 if (size() == 1) return V{1};
                                                                 } while (check(w) != P - 1);
                                                                  return qpow(S(P, w, a, 1), (P + 1) / 2).x;
 Poly X = Poly(*this, (size() + 1) / 2).Exp().isz(size
    ());
  Poly Y = X.Ln(); Y[0] = mod - 1;
                                                               5.23 Simplex
 fi(0, size()) Y[i] = modsub((*this)[i], Y[i]);
  return X.Mul(Y).isz(size());
                                                               namespace simplex {
                                                               // maximize c^Tx under Ax <= B
Poly Pow(const string &K) const {
                                                               // return VD(n, -inf) if the solution doesn't exist
 int nz = 0;
                                                               // return VD(n, +inf) if the solution is unbounded
  while (nz < size() && !(*this)[nz]) ++nz;</pre>
                                                               using VD = vector<double>;
  int nk = 0, nk2 = 0;
                                                               using VVD = vector<vector<double>>;
                                                               const double eps = 1e-9;
  for (char c : K) {
  nk = (nk * 10 + c - '0') % mod;
                                                               const double inf = 1e+9;
   nk2 = nk2 * 10 + c - '0';
                                                               int n, m;
  if (nk2 * nz >= size())
                                                               VVD d:
    return Poly(size());
                                                               vector<int> p, q;
   nk2 %= mod - 1;
                                                               void pivot(int r, int s) {
                                                                double inv = 1.0 / d[r][s];
  if (!nk && !nk2) return Poly(V{1}, size());
                                                                for (int i = 0; i < m + 2; ++i)
  Poly X = V(data() + nz, data() + size() - nz * (nk2 -
                                                                 for (int j = 0; j < n + 2; ++j)
     1));
                                                                   if (i != r && j != s)
  int x0 = X[0];
                                                                   d[i][j] = d[r][j] * d[i][s] * inv;
  return X.imul(modinv(x0)).Ln().imul(nk).Exp().imul(
                                                                for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
    modpow(x0, nk2)).irev().isz(size()).irev();
                                                                d[r][s] = inv; swap(p[r], q[s]);
Poly InvMod(int L) { // (to evaluate linear recursion)
 Poly R{1, 0}; // *this * R mod x^L = 1 (*this[0] ==
                                                               bool phase(int z) {
    1)
                                                                int x = m + z;
  for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                                while (true) {
   Poly 0 = R.Mul(Poly(data(), min<int>(2 << level,
                                                                 int s = -1:
                                                                 for (int i = 0; i <= n; ++i) {
  if (!z && q[i] == -1) continue</pre>
    size())))
   Poly Q(2 << level); Q[0] = 1;
   for (int j = (1 << level); j < (2 << level); ++j)</pre>
                                                                  if (s == -1) \mid d[x][i] < d[x][s]) s = i;
   Q[j] = modsub(mod, O[j]);
   R = R.Mul(Q).isz(4 << level);
                                                                 if (d[x][s] > -eps) return true;
  }
                                                                 int r = -1;
                                                                 for (int i = 0; i < m; ++i) {</pre>
  return R.isz(L);
                                                                  if (d[i][s] < eps) continue;</pre>
                                                                  if (r == -1 ||
static int LinearRecursion(const V &a, const V &c,
    int64_t n) { // a_n = \sum c_j a_(n-j)}
                                                                    d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  const int k = (int)a.size();
                                                                 if (r == -1) return false;
  assert((int)c.size() == k + 1);
  Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
                                                                 pivot(r, s);
 fi(1, k + 1) C[k - i] = modsub(mod, c[i]);
 C[k] = 1
  while (n) {
                                                               VD solve(const VVD &a, const VD &b, const VD &c) {
  if (n % 2) W = W.Mul(M).DivMod(C).second;
                                                                m = b.size(), n = c.size();
   n /= 2, M = M.Mul(M).DivMod(C).second;
                                                                d = VVD(m + 2, VD(n + 2));
                                                                for (int i = 0; i < m; ++i)</pre>
  int ret = 0;
                                                                 for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
  fi(0, k) ret = modadd(ret, modmul(W[i], a[i]));
                                                                p.resize(m), q.resize(n + 1);
                                                                for (int i = 0; i < m; ++i)
  return ret:
                                                                 p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
                                                                for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];
q[n] = -1, d[m + 1][n] = 1;
#undef fi
using Poly_t = Poly<998244353, 3, 1 << 20>;
                                                                int r = 0;
                                                                for (int i = 1; i < m; ++i)
if (d[i][n + 1] < d[r][n + 1]) r = i;</pre>
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
5.22 Quadratic residue
                                                                if (d[r][n + 1] < -eps) {</pre>
struct S {
                                                                 pivot(r, n);
int MOD, w;
                                                                 if (!phase(1) \mid | d[m + 1][n + 1] < -eps)
int64_t x, y;
                                                                  return VD(n, -inf);
                                                                 for (int i = 0; i < m; ++i) if (p[i] == -1) {
S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
  : MOD(m), w(w_{-}), x(x_{-}), y(y_{-}) {}
                                                                  int s = min_element(d[i].begin(), d[i].end() - 1)
S operator*(const S &rhs) const {
                                                                       - d[i].begin();
```

template <typename U, typename V>

```
pivot(i, s);
                                                                 bool isInter(U A, V B) {
                                                                   if (cross(A.dir, B.dir) == 0)
                                                                    return // handle parallel yourself
                                                                     isInter(A, B.st) || isInter(A, B.st+B.dir) ||
isInter(B, A.st) || isInter(B, A.st+A.dir);
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)
                                                                   Point D = B.st - A.st;
 if (p[i] < n) x[p[i]] = d[i][n + 1];
                                                                   11d C = cross(A.dir, B.dir)
                                                                   return U::valid(cross(D, A.dir), C) &&
 return x;
                                                                     V::valid(cross(D, B.dir), C);
5.24 Simplex Construction
                                                                  struct Line {
                                                                   Point st, ed, dir;
Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that for all 1 \leq j \leq m,
                                                                   Line (Point s, Point e)
\sum_{1 \le i \le n} A_{ji} x_i \le b_j and x_i \ge 0 for all 1 \le i \le n.
                                                                    : st(s), ed(e), dir(e - s) {}
  1. In case of minimization, let c_i^\prime = -c_i
                                                                 Pointf intersect(const Line &A, const Line &B) {
  2. \sum_{1 < i < n} A_{ji} x_i \ge b_j \to \sum_{1 < i < n} -A_{ji} x_i \le -b_j
                                                                   11f t = cross(B.st - A.st, B.dir) /
                                                                    llf(cross(A.dir, B.dir));
  3. \sum_{1 < i < n} A_{ji} x_i = b_j
                                                                   return toPointf(A.st) +
                                                                    Pointf(t) * toPointf(A.dir);
        • \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
        • \sum_{1 \le i \le n} A_{ji} x_i \ge b_j
                                                                  6.3 2D Convex Hull
                                                                 template<typename PT>
  4. If x_i has no lower bound, replace x_i with x_i - x_i'
                                                                  vector<PT> buildConvexHull(vector<PT> d) {
                                                                   sort(ALL(d), [](const PT& a, const PT& b){
     Geometry
                                                                     return tie(a.x, a.y) < tie(b.x, b.y);});</pre>
                                                                   vector<PT> s(SZ(d)<<1);</pre>
     Basic Geometry
                                                                   int o = 0;
#define IM imag
                                                                   for(auto p: d) {
#define RE real
                                                                    while(o \ge 2 \& cross(p-s[o-2], s[o-1]-s[o-2]) <= 0)
using lld = int64_t;
using llf = long double;
                                                                    s[o++] = p;
using Point = std::complex<11d>;
using Pointf = std::complex<llf>;
                                                                   for(int i=SZ(d)-2, t = o+1; i>=0; i--){
auto toPointf(Point p) { return Pointf(IM(p), RE(p)); }
                                                                    while(o>=t&&cross(d[i]-s[o-2],s[o-1]-s[o-2])<=0)
int sgn(11d x) \{ return (x > 0) - (x < 0); \}
                                                                     0--
lld dot(Point a, Point b) { return RE(conj(a) * b); }
lld cross(Point a, Point b) { return IM(conj(a) * b); }
                                                                    s[o++] = d[i];
int ori(Point a, Point b, Point c) {
                                                                   s.resize(o-1);
return sgn(cross(b - a, c - a));
                                                                   return s;
bool operator<(const Point &a, const Point &b) {</pre>
                                                                  6.4 3D Convex Hull
return RE(a) != RE(b) ? RE(a) < RE(b) : IM(a) < IM(b);
                                                                  // return the faces with pt indexes
int argCmp(Point a, Point b) {
                                                                  int flag[MXN][MXN];
 // -1 / 0 / 1 <-> < / == / > (atan2)
                                                                  struct Point{
 int qa = (IM(a) == 0
                                                                   ld x,y,z;
   ? (RE(a) < 0 ? 3 : 1) : (IM(a) < 0 ? 0 : 2));
                                                                   Point operator * (const 1d &b) const {
 int qb = (IM(b) == 0
                                                                    return (Point) {x*b, y*b, z*b};}
   ? (RE(b) < 0 ? 3 : 1) : (IM(b) < 0 ? 0 : 2));
                                                                   Point operator * (const Point &b) const {
 if (qa != qb)
                                                                    return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
  return sgn(qa - qb);
 return sgn(cross(b, a));
                                                                 Point ver(Point a, Point b, Point c) {
  return (b - a) * (c - a);}
template <typename V> llf area(const V & pt) {
 11d ret = 0;
                                                                  vector<Face> convex_hull_3D(const vector<Point> pt) {
 for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
                                                                   int n = SZ(pt), ftop = 0;
  ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
                                                                   REP(i,n) REP(j,n) flag[i][j] = 0;
 return ret / 2.0;
                                                                   vector<Face> now;
                                                                   now.emplace_back(0,1,2);
Point rot90(Point p) { return Point{-IM(p), RE(p)}; }
                                                                   now.emplace_back(2,1,0)
Pointf project(Pointf p, Pointf q) { // p onto q
                                                                   for (int i=3; i<n; i++){</pre>
return dot(p, q) * q / dot(q, q);
                                                                    ftop++; vector<Face> next;
                                                                    REP(j, SZ(now)) {
  Face& f=now[j]; int ff = 0;
6.2 Segment & Line Intersection
                                                                     ld d=(pt[i]-pt[f.a]).dot(
                                                                       ver(pt[f.a], pt[f.b], pt[f.c]));
struct Segment {
 Point st, dir; // represent st + t*dir for 0<=t<=1
                                                                     if (d <= 0) next.push_back(f);</pre>
 Segment(Point s, Point e) : st(s), dir(e - s) {}
                                                                     if (d > 0) ff=ftop;
 static bool valid(lld p, lld q) {
                                                                     else if (d < 0) ff=-ftop;</pre>
                                                                     flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
  // is there t s.t. 0 <= t <= 1 && qt == p ?
  if (q < 0) q = -q, p = -p;
                                                                    REP(j, SZ(now)) {
  Face& f=now[j];
  return 0 <= p && p <= q;
                                                                     if (flag[f.a][f.b] > 0 &&
   flag[f.a][f.b] != flag[f.b][f.a])
bool isInter(Segment A, Point P) {
 if (A.dir == Point(0)) return P == A.st;
                                                                      next.emplace_back(f.a,f.b,i);
 return cross(P - A.st, A.dir) == 0 &&
                                                                     if (flag[f.b][f.c] > 0 &&
  Segment::valid(dot(P - A.st, A.dir), norm(A.dir));
                                                                        flag[f.b][f.c] != flag[f.c][f.b])
```

next.emplace_back(f.b,f.c,i);

if (flag[f.c][f.a] > 0 &&

```
flag[f.c][f.a] != flag[f.a][f.c])
    next.emplace_back(f.c,f.a,i);
                                                                  return d;
 now=next;
                                                                 6.8 Simulated Annealing
 return now;
                                                                11f anneal() {
                                                                  mt19937 rnd_engine( seed );
                                                                  uniform_real_distribution< llf > rnd( 0, 1 );
6.5 2D Farthest Pair
                                                                  const 11f dT = 0.001;
// stk is from convex hull
                                                                  // Argument p
                                                                  11f S_cur = calc( p ), S_best = S_cur;
for ( 11f T = 2000 ; T > EPS ; T -= dT ) {
n = (int)(stk.size());
int pos = 1, ans = 0; stk.push_back(stk[0]);
for(int i=0;i<n;i++) {</pre>
                                                                   // Modify p to p_prime
                                                                   const llf S_prime = calc( p_prime );
 while(abs(cross(stk[i+1]-stk[i],
                                                                   const llf delta_c = S_prime - S_cur;
llf prob = min( ( llf ) 1, exp( -delta_c / T ) );
   stk[(pos+1)%n]-stk[i])) >
   abs(cross(stk[i+1]-stk[i],
                                                                   if ( rnd( rnd_engine ) <= prob )</pre>
   stk[pos]-stk[i]))) pos = (pos+1)%n;
                                                                   S_cur = S_prime, p = p_prime;
if ( S_prime < S_best ) // find min</pre>
 ans = max({ans, dis(stk[i], stk[pos]),
  dis(stk[i+1], stk[pos])});
                                                                    S_best = S_prime, p_best = p_prime;
6.6 2D Closest Pair
                                                                  return S_best;
struct cmp_y {
 bool operator()(const P& p, const P& q) const {
                                                                      Half Plane Intersection
  return p.y < q.y;</pre>
                                                                 // cross(pt-line.st, line.dir)<=0 <-> pt in half plane
                                                                bool operator<(const Line &lhs, const Line &rhs) {</pre>
multiset<P, cmp_y> s;
                                                                   if (int cmp = argCmp(lhs.dir, rhs.dir))
void solve(P a[], int n) {
                                                                     return cmp == -1;
                                                                   return ori(lhs.st, lhs.ed, rhs.st) < \theta;
 sort(a, a + n, [](const P& p, const P& q) {
  return tie(p.x, p.y) < tie(q.x, q.y);</pre>
 11f d = INF; int pt = 0;
for (int i = 0; i < n; ++i) {</pre>
                                                                 // intersect function is in "Segment Intersect"
                                                                llf HPI(vector<Line> &lines) {
                                                                   sort(lines.begin(), lines.end());
  while (pt < i \text{ and } a[i].x - a[pt].x >= d)
   s.erase(s.find(a[pt++]));
                                                                   deque<Line> que;
                                                                   deque<Pointf> pt;
  auto it = s.lower_bound(P(a[i].x, a[i].y - d));
                                                                   que.push_back(lines[0]);
  while (it != s.end() and it->y - a[i].y < d)
   d = min(d, dis(*(it++), a[i]));
                                                                   for (int i = 1; i < (int)lines.size(); i++) {</pre>
                                                                     if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
  s.insert(a[i]);
                                                                      continue;
}
                                                                 #define POP(L, R) \
                                                                     while (pt.size() > 0 \
     kD Closest Pair (3D ver.)
                                                                       && ori(L.st, L.ed, pt.back()) < 0) \
                                                                     pt.pop_back(), que.pop_back(); \
while (pt.size() > 0 \
11f solve(vector<P> v) {
 shuffle(v.begin(), v.end(), mt19937());
                                                                       && ori(R.st, R.ed, pt.front()) < 0) \
 unordered_map<1ld, unordered_map<1ld,</pre>
                                                                     pt.pop_front(), que.pop_front();
POP(lines[i], lines[i]);
  unordered_map<lld, int>>> m;
 llf d = dis(v[0], v[1]);
                                                                     pt.push_back(intersect(que.back(), lines[i]));
 auto Idx = [\&d] (11f x) -> 11d {
                                                                     que.push_back(lines[i]);
  return round(x * 2 / d) + 0.1;
 auto rebuild_m = [&m, &v, &Idx](int k) {
                                                                   POP(que.front(), que.back())
  m.clear();
                                                                   if (que.size() <= 1 ||</pre>
  for (int i = 0; i < k; ++i)
m[Idx(v[i].x)][Idx(v[i].y)]</pre>
                                                                     argCmp(que.front().dir, que.back().dir) == 0)
    [Idx(v[i].z)] = i;
                                                                   pt.push_back(intersect(que.front(), que.back()));
 }; rebuild_m(2);
                                                                   return area(pt);
 for (size_t i = 2; i < v.size(); ++i) {</pre>
  const 11d kx = Idx(v[i].x), ky = Idx(v[i].y),
  kz = Idx(v[i].z); bool found = false;
for (int dx = -2; dx <= 2; ++dx) {
                                                                 6.10 Minkowski Sum
                                                                 vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
   const 11d nx = dx + kx;
                                                                  hull(A), hull(B);
   if (m.find(nx) == m.end()) continue;
                                                                  vector<pll> C(1, A[0] + B[0]), s1, s2;
   auto& mm = m[nx];
                                                                  for(int i = 0; i < SZ(A); ++i)
s1.pb(A[(i + 1) % SZ(A)] - A[i]);
   for (int dy = -2; dy <= 2; ++dy) {
    const 11d ny = dy + ky;
                                                                  for(int i = 0; i < SZ(B); i++)</pre>
    if (mm.find(ny) == mm.end()) continue;
                                                                   s2.pb(B[(i + 1) % SZ(B)] - B[i])
    auto& mmm = mm[ny];
                                                                  for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
    for (int dz = -2; dz <= 2; ++dz) {
                                                                   if (p2 >= SZ(B)
     const 11d nz = dz + kz;
                                                                     || (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
     if (mmm.find(nz) == mmm.end()) continue;
                                                                    C.pb(C.back() + s1[p1++]);
     const int p = mmm[nz];
     if (dis(v[p], v[i]) < d) {
  d = dis(v[p], v[i]);</pre>
                                                                    C.pb(C.back() + s2[p2++]);
                                                                  return hull(C), C;
      found = true;
     }
                                                                 6.11 Circle Class
                                                                struct Circle { Pointf o; llf r; };
  if (found) rebuild_m(i + 1);
  else m[kx][ky][kz] = i;
                                                                vector<llf> intersectAngle(Circle A, Circle B) {
```

```
Pointf dir = B.o - A.o; llf d2 = norm(dir);
if (norm(A.r - B.r) >= d2)
 if (A.r < B.r) return {-PI, PI}; // special</pre>
else return {};
if (norm(A.r + B.r) <= d2) return {};</pre>
 11f dis = abs(dir), theta = arg(dir);
11f phi = acos((A.r * A.r + d2 - B.r * B.r) /
   (2 * A.r * dis));
11f L = theta - phi, R = theta + phi;
while (L < -PI) L += PI * 2;
while (R > PI) R -= PI * 2;
return { L, R };
vector<Pointf> intersectPoint(Circle a, Circle b) {
llf d = abs(a.o - b.o);
 if (d >= b.r+a.r || d <= abs(b.r-a.r)) return {};</pre>
11f dt = (b.r*b.r - a.r*a.r)/d, d1 = (d+dt)/2;
Pointf dir = (a.o - b.o); dir /= d;
Pointf pcrs = dir*d1 + b.o;
dt=sqrt(max(0.0L, b.r*b.r-d1*d1)), dir = rot90(dir);
return {pcrs + dir*dt, pcrs - dir*dt};
```

6.12 Intersection of line and Circle

```
vector<pdd> line_interCircle(const pdd &p1,
   const pdd &p2, const pdd &c, const double r) {
  pdd ft = p1 + project(c-p1, p2-p1), vec = p2-p1;
  llf dis = abs(c - ft);
  if (abs(dis - r) < eps) return {ft};
  if (dis > r) return {};
  vec = vec * sqrt(r * r - dis * dis) / abs(vec);
  return {ft + vec, ft - vec};
}
```

6.13 Intersection of Polygon and Circle

```
// Divides into multiple triangle, and sum up
// test by HDU2892
const double PI=acos(-1);
double _area(pdd pa, pdd pb, double r){
if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
if(abs(pb)<eps) return 0;</pre>
double S, h, theta;
double a=abs(pb), b=abs(pa), c=abs(pb-pa);
double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
double cosC = dot(pa,pb) / a / b, C = acos(cosC);
if(a > r){
 S = (C/2)*r*r
 h = a*b*sin(C)/c;
 if (h < r && B < PI/2)
  S = (acos(h/r)*r*r - h*sqrt(r*r-h*h));
else if(b > r){
 theta = PI - B - asin(sin(B)/r*a);
 S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
else S = .5*sin(C)*a*b;
return S;
double area_poly_circle(const vector<pdd> &poly,
 const pdd &0,const double r){
 double S=0; int N=poly.size();
for(int i=0;i<N;++i)</pre>
 S += _area(poly[i]-0, poly[(i+1)%N]-0, r)
* ori(0, poly[i], poly[(i+1)%N]);
return fabs(S);
```

6.14 Point & Hulls Tangent

```
#define above(P, Vi, Vj) (ori(P, Vi, Vj) > 0) // true
    if Vi is above Vj
#define below(P, Vi, Vj) (ori(P, Vi, Vj) < 0) // true
    if Vi is below Vj
// Rtangent_PointPolyC(): binary search for convex
    polygon right tangent
// Input: P = a 2D point (exterior to the polygon)
// n = number of polygon vertices
// V = array of vertices for a 2D convex polygon
    with V[n] = V[0]
// Return: index "i" of rightmost tangent point V[i]
int Rtangent_PointPolyC(Point P, int n, Point *V) {</pre>
```

```
int a, b, c
 int upA, dnC;
 if (below(P, V[1], V[0]) && !above(P, V[n - 1], V[0]))
  return 0:
 for (a = 0, b = n;;) {
  c = (a + b) / 2;
  dnC = below(P, V[c + 1], V[c]);
  if (dnC && !above(P, V[c - 1], V[c]))
   return c;
  upA = above(P, V[a + 1], V[a]);
  if (upA) {
   if (dnC) {
    b = c;
   } else {
    if (above(P, V[a], V[c]))
     b = c;
    else
     a = c:
  } else {
   if (!dnC) {
    a = c;
   } else {
    if (below(P, V[a], V[c]))
     b = c;
    else
     a = c;
}
// Ltangent_PointPolyC(): binary search for convex
    polygon left tangent
    Input: P = a 2D point (exterior to the polygon)
//
//
        n = number of polygon vertices
        V = array of vertices for a 2D convex polygon
//
    with V[n]=V[0]
    Return: index "i" of leftmost tangent point V[i]
int Ltangent_PointPolyC(Point P, int n, Point *V) {
 int a, b, c;
 int dnA, dnC;
 if (above(P, V[n - 1], V[0]) && !below(P, V[1], V[0]))
  return 0;
 for (a = 0, b = n;;) {
  c = (a + b) / 2;
  dnC = below(P, V[c + 1], V[c]);
  if (above(P, V[c - 1], V[c]) && !dnC)
  dnA = below(P, V[a + 1], V[a]);
  if (dnA) {
   if (!dnC) {
    b = c;
   } else {
    if (below(P, V[a], V[c]))
     b = c;
    else
  } else {
   if (dnC) {
    a = c;
   } else {
    if (above(P, V[a], V[c]))
     b = c;
    else
     a = c;
}
```

6.15 Convex Hulls Tangent

```
// RLtangent_PolyPolyC(): get the RL tangent between
   two convex polygons
// Input: m = number of vertices in polygon 1
```

```
V = array of vertices for convex polygon 1 with
                                                                  if(dist(pts[j], c.o) <= c.r) continue;</pre>
                                                                  c.o = (pts[i] + pts[j]) / 2;
     V[m]=V[0]
                                                                  c.r = dist(pts[i], c.o);
        n = number of vertices in polygon 2
//
                                                                 for (int k = 0; k < j; k++) {
  if (dist(pts[k], c.o) <= c.r) continue;</pre>
        W = array of vertices for convex polygon 2 with
     W[n]=W[0]
   Output: *t1 = index of tangent point V[t1] for
                                                                   c = getCircum(pts[i], pts[j], pts[k]);
    polygon 1
        *t2 = index of tangent point W[t2] for polygon
//
                                                                 }
void RLtangent_PolyPolyC(int m, Point *V, int n, Point
                                                                return c;
    *W, int *t1, int *t2) {
int ix1, ix2; // search indices for polygons 1 and 2
                                                                     KDTree (Nearest Point)
                                                               6.18
 // first get the initial vertex on each polygon
                                                              const int MXN = 100005;
ix1 = Rtangent_PointPolyC(W[0], m, V); // right
                                                              struct KDTree {
    tangent from W[0] to V
                                                                struct Node {
ix2 = Ltangent_PointPolyC(V[ix1], n, W); // left
                                                                 int x,y,x1,y1,x2,y2;
                                                                int id,f;
    tangent from V[ix1] to W
                                                                Node *L, *R;
                                                                } tree[MXN], *root;
 // ping-pong linear search until it stabilizes
int done = false; // flag when done
                                                                int n:
                                                               LL dis2(int x1, int y1, int x2, int y2) {
LL dx = x1-x2, dy = y1-y2;
while (done == false) {
 done = true; // assume done until..
 while (ori(W[ix2], V[ix1], V[ix1 + 1]) <= 0) {</pre>
                                                                return dx*dx+dy*dy;
  ++ix1; // get Rtangent from W[ix2] to V
                                                               static bool cmpx(Node& a, Node& b){return a.x<b.x;}
static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 while (ori(V[ix1], W[ix2], W[ix2 - 1]) >= 0) {
             // get Ltangent from V[ix1] to W
                                                                void init(vector<pair<int,int>> ip) {
   done = false; // not done if had to adjust this
                                                                n = ip.size();
  }
                                                                 for (int i=0; i<n; i++) {
                                                                 tree[i].id = i;
 *t1 = ix1;
                                                                  tree[i].x = ip[i].first;
*t2 = ix2;
                                                                  tree[i].y = ip[i].second;
return;
                                                                 root = build_tree(0, n-1, 0);
      Tangent line of Two Circle
6.16
                                                                Node* build_tree(int L, int R, int d) {
vector<Line> go(const Cir &c1, const Cir &c2,
                                                                 if (L>R) return nullptr;
 int sign1) {
// sign1 = 1 for outer tang, -1 for inter tang
                                                                 int M = (L+R)/2; tree[M].f = d%2;
                                                                 nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
vector<Line> ret;
                                                                 tree[M].x1 = tree[M].x2 = tree[M].x;
if (norm(c1.o - c2.o) < eps)
                                                                 tree[M].y1 = tree[M].y2 = tree[M].y;
                                                                 tree[M].L = build_tree(L, M-1, d+1);
  return ret;
11f d = abs(c1.o - c2.o);
                                                                 if (tree[M].L) {
                                                                 tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
Pointf v = (c2.o - c1.o) / d;
11f c = (c1.r - sign1 * c2.r) / d;
if (c * c > 1)
                                                                  tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
 return ret;
                                                                 tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
llf h = sqrt(max(0.0, 1.0 - c * c));
for (int sign2: {1, -1}) {
                                                                 tree[M].R = build_tree(M+1, R, d+1);
 Pointf n = c * v + sign2 * h * rot90(v);
                                                                 if (tree[M].R) {
                                                                 tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
 Pointf p1 = c1.o + n * c1.r;
 Pointf p2 = c2.0 + n * (c2.r * sign1);
 if (norm(p2 - p1) < eps)
                                                                  tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
  p2 = p1 + rot90(c2.o - c1.o);
                                                                 tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
 ret.push_back({p1, p2});
                                                                 return tree+M;
return ret;
                                                                int touch(Node* r, int x, int y, LL d2){
                                                                LL dis = sqrt(d2)+1;
     Minimum Covering Circle
                                                                 if (x<r->x1-dis || x>r->x2+dis ||
template<typename P>
                                                                  y<r->y1-dis || y>r->y2+dis)
Circle getCircum(const P &a, const P &b, const P &c){
                                                                  return 0;
Real a1 = a.x-b.x, b1 = a.y-b.y;
                                                                 return 1;
Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
Real a2 = a.x-c.x, b2 = a.y-c.y;
                                                                void nearest(Node* r,int x,int y,int &mID,LL &md2) {
                                                                if (!r || !touch(r, x, y, md2)) return;
LL d2 = dis2(r->x, r->y, x, y);
Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
Circle cc;
                                                                 if (d2 < md2 || (d2 == md2 && mID < r->id)) {
cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
                                                                 mID = r -> id;
cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);
                                                                 md2 = d2;
return cc;
                                                                 // search order depends on split dim
                                                                 if ((r->f == 0 \&\& x < r->x) ||
                                                                   (r->f == 1 && y < r->y)) {
template<typename P>
Circle MinCircleCover(const vector<P>& pts){
                                                                  nearest(r->L, x, y, mID, md2);
                                                                 nearest(r->R, x, y, mID, md2);
random_shuffle(pts.begin(), pts.end());
Circle c = { pts[0], 0 };
                                                                 } else {
for(int i=0;i<(int)pts.size();i++){</pre>
                                                                 nearest(r->R, x, y, mID, md2);
 if (dist(pts[i], c.o) <= c.r) continue;</pre>
                                                                  nearest(r->L, x, y, mID, md2);
  c = { pts[i], 0 };
 for (int j = 0; j < i; j++) {
```

```
int query(int x, int y) {
                                                                      11f B = atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);
  int id = 1029384756;
                                                                     eve[E++] = Teve(bb,B,1), eve[E++] = Teve(aa,A,-1);
  LL d2 = 102938475612345678LL;
                                                                     if(B > A) ++cnt;
  nearest(root, x, y, id, d2);
                                                                   if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
 return id:
}
                                                                   else{
} tree;
                                                                    sort(eve, eve + E);
                                                                    eve[E] = eve[0];
6.19
      Rotating Sweep Line
                                                                     for(int j = 0; j < E; ++j){
void rotatingSweepLine(pair<int, int> a[], int n) {
                                                                     cnt += eve[j].add;
                                                                     Area[cnt] += cross(eve[j].p, eve[j + 1].p) * .5;
 vector<pair<int, int>> 1;
 1.reserve(n * (n - 1) / 2);
                                                                      double theta = eve[j + 1].ang - eve[j].ang;
 for (int i = 0; i < n; ++i)
                                                                      if (theta < 0) theta += 2. * pi;</pre>
  for (int j = i + 1; j < n; ++j)
                                                                      Area[cnt]+=(theta-sin(theta))*c[i].R*c[i].R*.5;
   1.emplace_back(i, j);
 sort(1.begin(), 1.end(), [&a](auto &u, auto &v){
  lld udx = a[u.first].first - a[u.second].first;
lld udy = a[u.first].second - a[u.second].second;
  1ld vdx = a[v.first].first - a[v.second].first;
  11d vdy = a[v.first].second - a[v.second].second;
  if (udx == 0 \text{ or } vdx == 0) \text{ return not } udx == 0;
                                                                     Stringology
  int s = sgn(udx * vdx);
  return udy * vdx * s < vdy * udx * s;
                                                                7.1 Hash
 });
                                                                class Hash {
 vector < int > idx(n), p(n);
 iota(idx.begin(), idx.end(), 0);
sort(idx.begin(), idx.end(), [&a](int i, int j){
                                                                  static constexpr int P = 127, Q = 1051762951;
                                                                  vector<int> h, p;
  return a[i] < a[j]; });
 for (int i = 0; i < n; ++i) p[idx[i]] = i;</pre>
                                                                  void init(const string &s){
 for (auto [i, j]: 1) {
                                                                   h.assign(s.size()+1, 0); p.resize(s.size()+1);
  // do here
                                                                   for (size_t i = 0; i < s.size(); ++i)</pre>
  swap(p[i], p[j]);
                                                                    h[i + 1] = add(mul(h[i], P), s[i]);
  idx[p[i]] = i, idx[p[j]] = j;
                                                                   generate(p.begin(), p.end(),[x=1,y=1,this]()
mutable{y=x;x=mul(x,P);return y;});
6.20
      Circle Cover
                                                                  int query(int 1, int r){ // 1-base (1, r]
                                                                   return sub(h[r], mul(h[1], p[r-1]));}
const int N = 1021;
                                                               };
struct CircleCover {
 int C
                                                                7.2 Suffix Array
 Cir c[N];
 bool g[N][N], overlap[N][N];
                                                                namespace sfx {
                                                                bool _t[maxn * 2];
 // Area[i] : area covered by at least i circles
 double Area[ N ];
                                                                int hi[maxn], rev[maxn];
                                                                int _s[maxn * 2], sa[maxn * 2], _c[maxn * 2];
int x[maxn], _p[maxn], _q[maxn * 2];
// sa[i]: sa[i]-th suffix is the
 void init(int _C){ C = _C;}
 struct Teve {
 pdd p; double ang; int add;
  Teve() {}
                                                                // i-th lexigraphically smallest suffix.
  Teve(pdd _a, double _b, int _c):p(_a), ang(_b), add(
                                                                // hi[i]: longest common prefix
                                                                // of suffix sa[i] and suffix sa[i - 1].
  bool operator<(const Teve &a)const
                                                                void pre(int *a, int *c, int n, int z) {
  {return ang < a.ang;}
                                                                 memset(a, 0, sizeof(int) * n);
 }eve[N * 2];
                                                                 memcpy(x, c, sizeof(int) * z);
 // strict: x = 0, otherwise x = -1
 bool disjuct(Cir &a, Cir &b, int x)
{return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
                                                                void induce(int *a,int *c,int *s,bool *t,int n,int z){
                                                                 memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
 bool contain(Cir &a, Cir &b, int x)
 {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
                                                                  if (a[i] && !t[a[i] - 1])
 bool contain(int i, int j) {
                                                                   a[x[s[a[i] - 1]]++] = a[i] - 1;
                                                                 memcpy(x, c, sizeof(int) * z);
  /* c[j] is non-strictly in c[i]. */
                                                                 for (int i = n - 1; i >= 0; --i)
if (a[i] && t[a[i] - 1])
  return (sign(c[i].R - c[j].R) > 0 \mid \mid (sign(c[i].R - c[i].R) \mid c[i])
    [j].R) == 0 \&\& i < j)) \&\& contain(c[i], c[j], -1);
                                                                   a[--x[s[a[i] - 1]]] = a[i] - 1;
 void solve(){
  fill_n(Area, C + 2, 0);
                                                                void sais(int *s, int *a, int *p, int *q,
                                                                 bool *t, int *c, int n, int z) {
  for(int i = 0; i < C; ++i)</pre>
   for(int j = 0; j < C; ++j)
                                                                 bool uniq = t[n - 1] = true;
    overlap[i][j] = contain(i, j);
                                                                 int nn=0, nmxz=-1, *nsa = a+n, *ns=s+n, last=-1;
  for(int i = 0; i < C; ++i)
                                                                 memset(c, 0, sizeof(int) * z);
                                                                 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
   for(int j = 0; j < C; ++j)
                                                                 for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
    g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                                                                 if (uniq) {
      disjuct(c[i], c[j], -1));
  for(int i = 0; i < C; ++i){
                                                                  for (int i = 0; i < n; ++i) a[--c[s[i]]] = i;
   int E = 0, cnt = 1;
                                                                  return:
   for(int j = 0; j < C; ++j)</pre>
    if(j != i && overlap[j][i])
                                                                 for (int i = n - 2; i >= 0; --i)
     ++cnt;
                                                                  t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
                                                                 pre(a, c, n, z);
for (int i = 1; i <= n - 1; ++i)
   for(int j = 0; j < C; ++j)
    if(i != j && g[i][j]) {
                                                                  if (t[i] && !t[i - 1])
     pdd aa, bb;
     CCinter(c[i], c[j], aa, bb);
                                                                   a[--x[s[i]]] = p[q[i] = nn++] = i;
     11f A = atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);
                                                                 induce(a, c, s, t, n, z);
```

```
for (int i = 0; i < n; ++i)
  if (a[i] && t[a[i]] && !t[a[i] - 1]) {
 bool neq = last < 0 || \</pre>
  memcmp(s + a[i], s + last,
(p[q[a[i]] + 1] - a[i]) * sizeof(int));
                                                                int run(const char* s) {
                                                                 int now = root;
 ns[q[last = a[i]]] = nmxz += neq;
                                                                 for (char c; c = *s; ++s) {
                                                                   if (!st[now].ch[c -= 'a']) return 0;
sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
                                                                   now = st[now].ch[c];
pre(a, c, n, z);
for (int i = nn - 1; i >= 0; --i)
                                                                 return st[now].cnt;
 a[-x[s[p[nsa[i]]]] = p[nsa[i]];
                                                                }
                                                               } SAM;
 induce(a, c, s, t, n, z);
                                                               7.4 KMP
void build(const string &s) {
const int n = int(s.size());
                                                               vector<int> kmp(const string &s) {
for (int i = 0; i < n; ++i) _s[i] = s[i];
                                                                vector<int> f(s.size(), 0);
                                                                /* f[i] = length of the longest prefix
 _s[n] = 0; // s shouldn't contain 0
sais(_s, sa, _p, _q, _t, _c, n + 1, 256);
for(int i = 0; i < n; ++i) rev[sa[i] = sa[i+1]] = i;</pre>
                                                                   (excluding s[0:i]) such that it coincides
                                                                  with the suffix of s[0:i] of the same length */
 int ind = hi[0] = 0;
                                                                /* i + 1 - f[i] is the length of the
for (int i = 0; i < n; ++i) {
                                                                  smallest recurring period of s[0:i] */
 if (!rev[i]) {
                                                                int k = 0;
  ind = 0:
                                                                for (int i = 1; i < (int)s.size(); ++i) {</pre>
                                                                 while (k > 0 \&\& s[i] != s[k]) k = f[k-1];
  continue;
                                                                 if (s[i] == s[k]) ++k;
 while (i + ind < n && \
                                                                 f[i] = k;
  s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
 hi[rev[i]] = ind ? ind-- : 0;
                                                                return f;
                                                               }
                                                               vector<int> search(const string &s, const string &t) {
                                                                // return 0-indexed occurrence of t in s
     Suffix Automaton
                                                                vector<int> f = kmp(t), r;
for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
struct SuffixAutomaton {
                                                                 while(k > 0 && (k==(int)t.size() || s[i]!=t[k]))
struct node {
                                                                  k = f[k - 1]
  int ch[K], len, fail, cnt, indeg;
                                                                 if (s[i] == t[k]) ++k;
 node(int L = 0) : ch{}, len(L), fail(0), cnt(0),
                                                                 if (k == (int)t.size()) r.push_back(i-t.size()+1);
    indeg(0) \{ \}
 } st[N];
                                                                return res;
 int root, last, tot;
                                                               }
void extend(int c) {
 int cur = ++tot;
                                                               7.5 Z value
 st[cur] = node(st[last].len + 1);
                                                               vector<int> Zalgo(const string &s) {
  while (last && !st[last].ch[c]) {
                                                                vector<int> z(s.size(), s.size());
    st[last].ch[c] = cur;
                                                                for (int i = 1, 1 = 0, r = 0; i < z[0]; ++i) {
  int j = clamp(r - i, 0, z[i - 1]);</pre>
    last = st[last].fail;
                                                                 for (; i + j < z[0] \text{ and } s[i + j] == s[j]; ++j);
  if (!last) {
    st[cur].fail = root;
                                                                 if (i + (z[i] = j) > r) r = i + z[1 = i];
  } else {
    int q = st[last].ch[c];
                                                                return z:
                                                               }
    if (st[q].len == st[last].len + 1) {
      st[cur].fail = q;
                                                               7.6 Manacher
    } else {
      int clone = ++tot;
                                                               int z[maxn];
      st[clone] = st[q];
st[clone].len = st[last].len + 1;
                                                               int manacher(const string& s) {
  string t = ".";
                                                                for(char c: s) t += c, t += '.';
      st[st[cur].fail = st[q].fail = clone].cnt = 0;
      while (last && st[last].ch[c] == q) {
                                                                int 1 = 0, r = 0, ans = 0;
        st[last].ch[c] = clone;
                                                                for (int i = 1; i < t.length(); ++i) {</pre>
                                                                 z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
        last = st[last].fail;
                                                                 while (i - z[i] >= 0 && i + z[i] < t.length()) {
      }
   }
                                                                  if(t[i - z[i]] == t[i + z[i]]) ++z[i];
                                                                  else break:
 st[last = cur].cnt += 1;
                                                                 if (i + z[i] > r) r = i + z[i], l = i;
void init(const char* s) {
 root = last = tot = 1;
                                                                for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);</pre>
  st[root] = node(0);
                                                                return ans;
  for (char c; c = *s; ++s) extend(c - 'a');
                                                               7.7 Lexico Smallest Rotation
int q[N];
void dp() {
                                                               string mcp(string s) {
  for (int i = 1; i <= tot; i++) ++st[st[i].fail].indeg</pre>
                                                                int n = s.length();
                                                                s += s; int i = 0, j = 1;
  int head = 0, tail = 0;
                                                                while (i < n && j < n) {</pre>
  for (int i = 1; i <= tot; i++)</pre>
                                                                 int k = 0;
    if (st[i].indeg == 0) q[tail++] = i;
                                                                 while (k < n \&\& s[i + k] == s[j + k]) k++;
  while (head != tail) {
                                                                 ((s[i+k] \le s[j+k])?j:i) += k+1;
    int now = q[head++];
                                                                 j += (i == j);
    if (int f = st[now].fail) {
   st[f].cnt += st[now].cnt;
                                                                return s.substr(i < n ? i : j, n);</pre>
      if (--st[f].indeg == 0) q[tail++] = f;
```

7.8 BWT

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a
 vector<int> v[ SIGMA ];
 void BWT(char* ori, char* res){
  // make ori -> ori + ori
  // then build suffix array
 void iBWT(char* ori, char* res){
  for( int i = 0 ; i < SIGMA ; i ++ )</pre>
   v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )</pre>
   v[ ori[i] - BASE ].push_back( i );
  vector<int> a;
  for(int i = 0,
                    ptr = 0 ; i < SIGMA ; i ++ )
   for( auto j : v[ i ] ){
    a.push_back( j );
    ori[ ptr ++ ] = BASE + i;
  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
   ptr = a[ ptr ];
  res[ len ] = 0;
} bwt;
```

7.9 Palindromic Tree

```
struct palindromic_tree{
struct node{
 int next[26],f,len;
 int cnt, num, st, ed;
 node(int l=0):f(0),len(1),cnt(0),num(0) {
  memset(next, 0, sizeof(next)); }
}:
vector<node> st;
vector<char> s;
int last, n;
void init(){
 st.clear();s.clear();last=1; n=0;
 st.push_back(0);st.push_back(-1);
 st[0].f=1;s.push_back(-1); }
int getFail(int x){
 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
 return x;}
void add(int c){
 s.push_back(c-='a'); ++n;
 int cur=getFail(last);
 if(!st[cur].next[c]){
  int now=st.size();
  st.push_back(st[cur].len+2);
  st[now].f=st[getFail(st[cur].f)].next[c];
  st[cur].next[c]=now;
  st[now].num=st[st[now].f].num+1;
 last=st[cur].next[c];
 ++st[last].cnt;}
 void dpcnt() {
 for (int i=st.size()-1; i >= 0; i--)
  st[st[i].f].cnt += st[i].cnt;
int size(){ return st.size()-2;}
} pt;
int main() {
string s; cin >> s; pt.init();
for (int i=0; i<SZ(s); i++) {</pre>
 int prvsz = pt.size(); pt.add(s[i]);
 if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [1,r]: s.substr(1, r-1+1)
 }
return 0;
```

8 Misc

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\dots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \le k \le n$.

8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W\subseteq X, |W|\leq |N_G(W)|$

8.1.7 Euler's planar graph formula

$$V - E + F = C + 1$$
, $E \le 3V - 6$ (?)

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Lucas's theorem

 $\binom{m}{n}\equiv\prod_{i=0}^k\binom{m_i}{n_i}\ (\mathrm{mod}\ p)\text{, where }m=m_kp^k+m_{k-1}p^{k-1}+\cdots+m_1p+m_0\text{,}$ and $n=n_kp^k+n_{k-1}p^{k-1}+\cdots+n_1p+n_0.$

8.1.10 Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2)$, find maximum $S\in I_1\cap I_2$. For each iteration, build the directed graph and find a shortest path from s to t.

- $s \to x : S \sqcup \{x\} \in I_1$
- $x \to t : S \sqcup \{x\} \in I_2$
- $y \to x : S \setminus \{y\} \sqcup \{x\} \in I_1$ (y is in the unique circuit of $S \sqcup \{x\}$)
- $x \to y: S \setminus \{y\} \sqcup \{x\} \in I_2$ (y is in the unique circuit of $S \sqcup \{x\}$)

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|.$ In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

8.2 DP-opt Condition

8.2.1 totally monotone (concave/convex)

$$\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}$$

8.2.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

if(!dfn[v]){

```
8.3 Convex 1D/1D DP
                                                                 par[v]=u;
                                                                 tarian(v):
struct segment {
                                                                 low[u]=min(low[u],low[v]);
 int i, l, r
                                                                 if(dfn[u]<low[v]){</pre>
 segment() {}
                                                                  g[u].push_back(v);
 segment(int \ a, \ int \ b, \ int \ c) \colon i(a), \ l(b), \ r(c) \ \{\}
                                                                  g[v].push_back(u);
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
                                                                }else{
void solve() {
                                                                 low[u]=min(low[u],dfn[v]);
 dp[0] = 0;
                                                                 if(dfn[v]<dfn[u]){</pre>
 deque<segment> dq; dq.push_back(segment(0, 1, n));
                                                                  int temp_v=u;
 for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i);</pre>
                                                                  bcc_id++;
                                                                  while(temp_v!=v){
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
                                                                   g[bcc_id+n].push_back(temp_v);
  dq.front().l = i + 1;
                                                                   g[temp_v].push_back(bcc_id+n);
  segment seg = segment(i, i + 1, n);
                                                                   temp_v=par[temp_v];
  while (dq.size() &&
   f(i, dq.back().1) < f(dq.back().i, dq.back().1))
                                                                  g[bcc_id+n].push_back(v);
    dq.pop_back();
                                                                  g[v].push_back(bcc_id+n);
  if (dq.size())
                                                                  reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
   int d = 1 << 20, c = dq.back().1;</pre>
   while (d >>= 1) if (c + d <= dq.back().r)</pre>
           c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
                                                             int dp[maxn][2], min_dp[2][2], tmp[2][2], tp[2];
  if (seg.1 <= n) dq.push_back(seg);</pre>
                                                              void dfs(int u,int fa){
                                                               if(u<=n){
                                                                for(int i=0;i<(int)g[u].size();i++){</pre>
                                                                 int v=g[u][i];
8.4
      ConvexHull Optimization
                                                                 if(v==fa) continue;
                                                                 dfs(v,u);
 mutable int64_t a, b, p;
                                                                 memset(tp,0x8f,sizeof tp);
 bool operator<(const L &r) const { return a < r.a; }</pre>
                                                                 if(v<=n){
 bool operator<(int64_t x) const { return p < x; }</pre>
                                                                  tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
                                                                  tp[1]=max(
struct DynamicHull : multiset<L, less<>> {
                                                                   dp[u][0]+dp[v][0]+1
 static const int64_t kInf = 1e18;
                                                                   dp[u][1]+max(dp[v][0],dp[v][1])
 bool Isect(iterator x, iterator y)
  auto Div = [](int64_t a, int64_t b) {
                                                                 }else{
    return a / b - ((a ^ b) < 0 && a % b); }
                                                                  tp[0]=dp[u][0]+dp[v][0];
  if (y == end()) { x->p = kInf; return false; }
                                                                  tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
  if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
                                                                 dp[u][0]=tp[0],dp[u][1]=tp[1];
  return x->p >= y->p;
                                                               }else{
 void Insert(int64_t a, int64_t b) {
                                                                for(int i=0;i<(int)g[u].size();i++){</pre>
  auto z = insert({a, b, 0}), y = z++, x = y;
                                                                 int v=g[u][i];
  while (Isect(y, z)) z = erase(z);
                                                                 if(v==fa) continue;
  if (x!=begin()\&Esct(--x,y)) Isect(x, y=erase(y));
                                                                 dfs(v,u);
  while ((y = x) != begin() && (--x)->p >= y->p)
   Isect(x, erase(y));
                                                                min_dp[0][0]=0;
                                                                min_dp[1][1]=1;
 int64_t Query(int64_t x) {
                                                                min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f3f;
 auto 1 = *lower_bound(x);
                                                                for(int i=0;i<(int)g[u].size();i++){</pre>
  return 1.a * x + 1.b;
                                                                 int v=g[u][i]
                                                                 if(v==fa) continue;
};
                                                                 memset(tmp,0x8f,sizeof tmp);
                                                                 tmp[0][0]=max(
8.5
      Josephus Problem
                                                                  min_dp[0][0]+max(dp[v][0],dp[v][1]),
// n people kill m for each turn
                                                                  min_dp[0][1]+dp[v][0]
int f(int n, int m) {
 int s = 0;
                                                                 tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
 for (int i = 2; i <= n; i++)
                                                                 tmp[1][0]=max(
  s = (s + m) \% i;
                                                                  \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
 return s;
                                                                  min_dp[1][1]+dp[v][0]
// died at kth
                                                                 tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
int kth(int n, int m, int k){
                                                                 memcpy(min_dp,tmp,sizeof tmp);
 if (m == 1) return n-1;
 for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
                                                                dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
 return k;
                                                                dp[u][0]=min_dp[0][0];
8.6 Cactus Matching
                                                              int main(){
vector<int> init_g[maxn],g[maxn*2];
                                                               int m,a,b;
                                                               scanf("%d%d",&n,&m);
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
                                                               for(int i=0;i<m;i++){
  scanf("%d%d",&a,&b);</pre>
void tarjan(int u){
 dfn[u]=low[u]=++dfs_idx;
 for(int i=0;i<(int)init_g[u].size();i++){</pre>
                                                                init_g[a].push_back(b);
  int v=init_g[u][i];
                                                                init_g[b].push_back(a);
  if(v==par[u]) continue;
```

par[1]=-1;

```
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 tarjan(1);
 dfs(1,-1);
 printf("%d\n", max(dp[1][0], dp[1][1]));
 return 0;
8.7
      Tree Knapsack
int dp[N][K]; PII obj[N];
vector<int> G[N];
void dfs(int u, int mx){
 for(int s: G[u]) {
  if(mx < obj[s].first) continue;</pre>
  for(int i=0;i<=mx-obj[s].FF;i++)</pre>
   dp[s][i] = dp[u][i];
  dfs(s, mx - obj[s].first);
  for(int i=obj[s].FF;i<=mx;i++)</pre>
   dp[u][i] = max(dp[u][i],
    dp[s][i - obj[s].FF] + obj[s].SS);
}
      N Queens Problem
8.8
vector< int > solve( int n ) {
 // no solution when n=2, 3
 vector< int > ret;
 if ( n % 6 == 2 ) {
 for ( int i = 2 ; i <= n ; i += 2 )
   ret.push_back( i );
 ret.push_back( 3 ); ret.push_back( 1 );
for ( int i = 7 ; i <= n ; i += 2 )
   ret.push_back( i );
  ret.push_back( 5 );
 } else if ( n % 6 == 3 ) {
  for ( int i = 4 ; i <= n ; i += 2 )
  ret.push_back( i );
  ret.push_back( 2 );
  for ( int i = 5 ; i <= n ; i += 2 )
  ret.push_back( i );
  ret.push_back( 1 ); ret.push_back( 3 );
 } else {
 for ( int i = 2 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
  for ( int i = 1 ; i \le n ; i += 2 )
   ret.push_back( i );
 return ret;
}
8.9
      Aliens Optimization
long long Alien() +
 long long c = kInf;
 for (int d = 60; d >= 0; --d) {
  // cost can be negative, depending on the problem.
  if (c - (1LL << d) < 0) continue;
  long long ck = c - (1LL \ll d);
  pair<long long, int> r = check(ck);
  if (r.second == k) return r.first - ck * k;
  if (r.second < k) c = ck;
 pair<long long, int> r = check(c);
 return r.first - c * k;
8.10
      Hilbert Curve
long long hilbert(int n, int x, int y) {
 long long res = 0;
 for (int s = n / 2; s; s >>= 1) {
 int rx = (x & s) > 0, ry = (y & s) > 0;
res += s * 111 * s * ((3 * rx) ^ ry);
  if (ry == 0) {
  if (rx == 1) x = s - 1 - x, y = s - 1 - y;
   swap(x, y);
  }
 return res;
```

```
8.11 Binary Search On Fraction
```

```
struct Q {
11 p, q;
 Q go(Q b, 11 d) \{ return \{p + b.p*d, q + b.q*d\}; \}
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(11 N) {
Q lo{0, 1}, hi{1, 0};
 if (pred(lo)) return lo;
 assert(pred(hi));
 bool dir = 1, L = 1, H = 1;
 for (; L || H; dir = !dir) {
 11 len = 0, step = 1;
  for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
   if (Q mid = hi.go(lo, len + step);
     mid.p > N || mid.q > N || dir ^ pred(mid))
    t++;
   else len += step;
  swap(lo, hi = hi.go(lo, len));
  (dir ? L : H) = !!len;
 return dir ? hi : lo;
```