Contents

1	Basi	ic										1
'	1.1	vimrc										1
	1.2	Debug Macro										1
	1.3	Increase Stack										i
	1.4	Pragma Optimization										2
	1.5	IO Optimization		 				 				2
_	_	_										_
2		a Structure										2
	2.1	Dark Magic										2
	2.2	Link-Cut Tree										2
	2.4	LiChao Segment Tree										3
	2.5	Sparse Table										3
	2.6	Linear Basis										3
	2.7	Binary Search On Segment Tree .										3
3	Grap											4
	3.1	BCC Edge										4
	3.2	BCC Vertex										4
	3.3 3.4	2-SAT (SCC)										4 5
	3.5	Lowbit Decomposition										5 5
	3.6	MaxCliqueDyn										6
	3.7	Virtural Tree										6
	3.8	Centroid Decomposition										6
	3.9	Tree Hashing		 				 				7
		Minimum Mean Cycle										7
	3.11	Mo's Algorithm on Tree										7
		Minimum Steiner Tree										7
		Directed Minimum Spanning Tree										8
		Dominator Tree										8
	5.15	Edge Coloring		 	٠	 ٠	•	 	٠	٠	 •	9
4	Mate	ching & Flow										9
•	4.1	Kuhn Munkres		 				 			 	9
	4.2	Bipartite Matching										9
	4.3	General Graph Matching										10
	4.4	Minimum Weight Matching (Clique										10
	4.5	Minimum Cost Circulation										10
	4.6	Flow Models										11
	4.7	Dinic										11
	4.8	Minimum Cost Maximum Flow										11
	4.9	GomoryHu Tree					•	 				12
		Global Min Cut										12
	4.10	Global Min-Cut		 			•	 				12
5	4.10 Mat		 •	 	•			 				12
5		h Prime Table		 				 			 	
5	Mat 5.1 5.2	h Prime Table $\dots \dots \dots$ $\lfloor rac{n}{i} \rfloor$ Enumeration $\dots \dots$		 				 			 	12 12 12
5	Mat 5.1 5.2 5.3	h Prime Table	 	 				 			 	12 12 12 12
5	Mat 5.1 5.2 5.3 5.4	h Prime Table	 	 		 		 			 	12 12 12 12 12
5	Mat 5.1 5.2 5.3 5.4 5.5	h Prime Table	 	 		 		 			 	12 12 12 12 12 13
5	Mat 5.1 5.2 5.3 5.4	h Prime Table	 	 		 		 			 	12 12 12 12 12 13
5	Mat 5.1 5.2 5.3 5.4 5.5	h Prime Table	 	 		 		 			 	12 12 12 12 12 13 13
5	Matl 5.1 5.2 5.3 5.4 5.5 5.6	h Prime Table	 	 		 						12 12 12 12 12 13 13 13
5	Mat 5.1 5.2 5.3 5.4 5.5	h Prime Table	 	 		 		 				12 12 12 12 12 13 13
5	Mattl 5.1 5.2 5.3 5.4 5.5 5.6	h Prime Table	 	 		 		 				12 12 12 12 12 13 13 13 13
5	Matl 5.1 5.2 5.3 5.4 5.5 5.6	h Prime Table	 			 						12 12 12 12 13 13 13 13 13
5	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11	h Prime Table										12 12 12 12 13 13 13 13 13 13 13 13
5	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12	h Prime Table										12 12 12 12 13 13 13 13 13 13 13 13
5	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13	h Prime Table										12 12 12 12 12 13 13 13 13 13 13 13 14 14
5	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14	h Prime Table $ \begin{bmatrix} \frac{n}{i} \end{bmatrix} $ Enumeration $3x+by=gcd$ Pollard Rho Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT										12 12 12 12 12 13 13 13 13 13 13 14 14 14
5	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15	h Prime Table $\begin{bmatrix} \frac{n}{i} \end{bmatrix}$ Enumeration $\frac{n}{i} \end{bmatrix}$ Enumeration $\frac{n}{i} \end{bmatrix}$ Enumeration $\frac{n}{i} \end{bmatrix}$ Polyneration $\frac{n}{i} \end{bmatrix}$ Polynomial Operations										12 12 12 12 12 13 13 13 13 13 13 13 14 14 14 14 15
5	Matl 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15	h Prime Table $\lfloor \frac{n}{i} \rfloor$ Enumeration $\frac{n}{i} \rfloor$ Enumeration $\frac{n}{i} \rfloor$ Enumeration $\frac{n}{i} \rfloor$ Pollard Rho Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT										12 12 12 12 13 13 13 13 13 13 13 14 14 14 15 16
5	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16	h Prime Table $\lfloor \frac{n}{4} \rfloor$ Enumeration $\frac{n}{4} \rfloor$ Enumeration $\frac{n}{4} \rfloor$ Enumeration $\frac{n}{4} \rfloor$ Pollard Rho Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog										12 12 12 12 12 13 13 13 13 13 13 14 14 14 14 15 16
5	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18	h Prime Table $\lfloor \frac{n}{i} \rfloor$ Enumeration $\frac{n}{i} \rfloor$ Enumeration $\frac{n}{i} \rfloor$ Enumeration $\frac{n}{i} \rfloor$ Pollard Rho Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT										12 12 12 12 13 13 13 13 13 13 13 14 14 14 15 16
5	Matt 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19	h Prime Table Prime Table \[\frac{\pi_1}{2} \] Enumeration \[\alpha \times \] Pollard Rho Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum										12 12 12 12 12 13 13 13 13 13 13 14 14 14 14 15 16 16
5	Matt 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20	h Prime Table Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum Quadratic residue										12 12 12 12 13 13 13 13 13 13 13 14 14 14 15 16 16 16
5	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.12 5.20 5.21 5.20	h Prime Table $\begin{bmatrix} \frac{n}{4} \end{bmatrix}$ Enumeration $\frac{n}{4}$ Full Enumeration Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum Quadratic residue De-Bruijn Simplex Construction Simplex										12 12 12 12 12 13 13 13 13 13 13 14 14 14 15 16 16 16 16 16 16 17 17
5	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.12 5.20 5.21 5.20	Prime Table										12 12 12 12 12 13 13 13 13 13 13 14 14 14 15 16 16 16 16 16
	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22 5.23	Prime Table Prime Table \[\frac{\pi}{2} \] Enumeration \[\frac{\pi}{2} \] Enumeration \[\frac{\pi}{2} \] Pollard Rho Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum Quadratic residue De-Bruijn Simplex Charateristic Polynomial										12 12 12 12 12 12 12 12 13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17
5	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22 5.23	Prime Table Prime Table 1/2 Enumeration										12 12 12 12 12 13 13 13 13 13 13 14 14 14 15 16 16 16 16 16 16 17 17
	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 5.22 5.23 Geo	Prime Table Prime Table \[\frac{\pi}{2} \] Enumeration \[\frac{\pi}{2} \] Enumeration \[\frac{\pi}{2} \] Pollard Rho Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum Quadratic residue De-Bruijn Simplex Charateristic Polynomial										12 12 12 12 12 12 12 12 13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17 18
	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.20 5.21 5.22 6.23 Geo 6.1	Prime Table Prime Table 1/2 Enumeration										12 12 12 12 12 12 13 13 13 13 13 13 13 14 14 15 16 16 16 17 17 17 18 18 18
	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.20 5.21 5.22 6.0 6.1 6.2	Prime Table Prime Table 1/2 Enumeration										12 12 12 12 12 12 13 13 13 13 13 13 13 14 14 15 16 16 16 16 17 17 17 17 18 18 18 18
	Matt 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 5.21 6.2 6.3	Prime Table Strling Number Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Charateristic Polynomial metry Basic Geometry Circle Class 2D Convex Hull 3D Convex Hull 2D Farthest Pair										122 122 122 122 133 133 133 133 133 134 144 155 166 166 177 177 177 188 188 188 188
	Mattl 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 6.17 5.20 6.2 6.3 6.4 6.5 6.6	Prime Table Prime Table Time Table Time Tab										12 12 12 12 12 12 13 13 13 13 13 13 13 14 14 15 16 16 16 17 17 17 18 18 18 18 18 18 18 18 18
	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 6.1 6.2 6.3 6.4 6.5 6.6 6.7	Prime Table Prime Table \[\frac{n}{2} \] Enumeration \[\frac{n}{2} \] Enumeration \[\frac{n}{2} \] Pollard Rho Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Charateristic Polynomial metry Basic Geometry Circle Class 2D Convex Hull 3D Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair kD Closest Pair kD Closest Pair kD Closest Pair (3D ver.)										12 12 12 12 12 12 13 13 13 13 13 13 13 13 14 14 15 16 16 16 17 17 17 17 18 18 18 18 18 18 18 19
	Matt 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 6.0 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8	Prime Table Prime Table Prime Table \[\frac{\pi}{2} \] Enumeration \[\ax+by=gcd \] Pollard Rho Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Charateristic Polynomial metry Basic Geometry Circle Class 2D Convex Hull 3D Convex Hull 3D Farthest Pair 2D Closest Pair kD Closest Pair kD Closest Pair kD Closest Pair Simulated Annealing										12 12 12 12 12 12 13 13 13 13 13 13 13 13 14 14 15 16 16 16 16 16 17 17 17 17 17 18 18 18 18 18 19 19 19
	Mattl 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.18 5.19 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9	Prime Table Prime Table Prime Table \[\frac{\pi}{2} \] Enumeration \[\alpha \text{x} + \begin{array}{c} \] Enumeration \[\alpha \text{x} + \begin{array}{c} \] \text{Pollard Rho} \[\text{Pi Count (Linear Sieve)} \] Strling Number \[\text{Sc1} \] Sirst Kind \[\text{5.6.2 Second Kind} \] Range Sieve \[\text{Miller Rabin} \] Extended Euler \[\text{Gauss Elimination} \] Fast Fourier Transform \[\text{Chinese Remainder} \] Berlekamp Massey \[\text{NTT} \] Polynomial Operations \[\text{FWT} \] DiscretLog FloorSum \[\text{Quadratic residue} \] \[\text{De-Bruijn} \] Simplex Construction \[\text{Simplex} \] Charateristic Polynomial \[\text{metry} \] Basic Geometry \[\text{Circle Class} \] \[\text{2D Convex Hull} \] \[\text{3D Convex Hull} \] \[\text{2D Closest Pair} \] \[\text{4D Closest Pair} \] \[4D Closes										122 122 122 133 133 133 133 133 134 144 155 166 166 167 177 177 178 188 188 188 189 199 199 199
	Mattl 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.17 5.20 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10	Prime Table Prime Table Table Ta										12 12 12 12 12 12 13 13 13 13 13 13 13 13 14 14 14 15 16 16 16 17 17 17 18 18 18 18 18 19 19 19 19 19
	Mattl 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.12 5.13 5.14 5.15 5.16 5.16 5.20 5.21 5.22 5.23 Geo 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11	Prime Table Prime Table ½ Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Charateristic Polynomial metry Basic Geometry Circle Class 2D Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair kD Closest Pair kD Closest Pair kD Closest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski sum intersection of line and circle										12 12 12 12 12 12 13 13 13 13 13 13 13 14 14 15 166 16 16 17 17 17 18 18 18 18 18 19 19 19 19 19 19 19
	Mati 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.23 5.21 5.22 5.23 Geo 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11 6.12	Prime Table Prime Table										12 12 12 12 12 12 13 13 13 13 13 13 13 13 14 14 15 16 16 16 17 17 17 18 18 18 18 19 19 19 19 19 19 19 19 19 19
	Mattl 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 5.11 5.12 5.13 5.14 5.15 5.16 5.17 5.18 6.0 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.12 6.13	Prime Table Prime Table ½ Enumeration ax+by=gcd Pollard Rho Pi Count (Linear Sieve) Strling Number 5.6.1 First Kind 5.6.2 Second Kind Range Sieve Miller Rabin Extended Euler Gauss Elimination Fast Fourier Transform Chinese Remainder Berlekamp Massey NTT Polynomial Operations FWT DiscreteLog FloorSum Quadratic residue De-Bruijn Simplex Construction Simplex Charateristic Polynomial metry Basic Geometry Circle Class 2D Convex Hull 3D Convex Hull 2D Farthest Pair 2D Closest Pair kD Closest Pair kD Closest Pair kD Closest Pair kD Closest Pair (3D ver.) Simulated Annealing Half Plane Intersection Minkowski sum intersection of line and circle										12 12 12 12 12 12 13 13 13 13 13 13 13 14 14 15 166 16 16 17 17 17 18 18 18 18 18 19 19 19 19 19 19 19

7	Strir	ngology	2
•	7.1	5 35	2
	7.2		2
	7.3		2
	7.4		22
	7.5		22
	7.5 7.6		22
	7.0 7.7		22
	7.7 7.8		22
	7.0 7.9		22
	7.9	Palinaromic free	
8	Misc	•	23
_	8.1		23
	•		23
			23
			23
			23
			23
			23
		· · · · · · · · · · · · · · · · · · ·	23
		- 9-	23
			23
			23
	8.2		23
	0.2		23
		3 · · · · · · · · · · · · · · · · · · ·	
	0.7		23
	8.3		23
	8.4	and the state of t	23
	8.5		23
	8.6	- · · · · · · · · · · · · · · · · · · ·	23
	8.7		24
	8.8		24
	8.9	~	25
	8.10		25
	8.11	To Check When Submit	25
1	_	Dania	
1		Basic	
1.	1	vimrc	
٠.	•	VIIII C	

```
se is nu bs=2 ru mouse=a encoding=utf-8 ls=2
se cin cino+=j1 et sw=4 sts=4 tgc sc hls
syn on
colorscheme desert
filetype indent on
inoremap {<CR> {<CR>}<ESC>0
map <F8> <ESC>:w<CR>:!g++ "%" -o "%<" -std=c++17 -
    DKISEKI -Wall -Wextra -Wshadow -Wfatal-errors -
    Wconversion -fsanitize=address,undefined -g && echo
     success<CR>
map <F9> <ESC>:w<CR>:!g++ "%" -o "%<" -02 -std=c++17 &&
     echo success<CR>
map <F10> <ESC>:!./"%<"<CR>
```

1.2 Debug Macro

```
#ifdef KISEKI
#define safe cerr<<__PRETTY_FUNCTION__\
<<" line "<<__LINE__<<" safe\n"</pre>
#define debug(a...) qwerty(#a, a)
#define orange(a...) dvorak(#a, a)
using std::cerr;
template <typename ...T>
void qwerty(const char *s, T ...a) {
  cerr << "\e[1;32m(" << s << ") = (";
  int cnt = sizeof...(T);</pre>
   (..., (cerr << a << (--cnt ? ", " : ")\e[0m\n")));
}
template <typename Iter>
void dvorak(const char *s, Iter L, Iter R) {
  cerr << "\e[1;32m[ " << s << " ] = [ ";</pre>
   for (int f = 0; L != R; ++L)
  cerr << (f++ ? ", " : "") << *L;
cerr << " ]\e[0m\n";</pre>
#else
#define safe ((void)0)
#define debug(...) ((void)\theta)
#define orange(...) ((void)\theta)
#endif
```

1.3 Increase Stack

```
const int size = 256 << 20;</pre>
                                    register long rsp asm("rsp");
                                    char *p = (char*)malloc(size)+size, *bak = (char*)rsp;
                                    __asm__("movq %0, %%rsp\n"::"r"(p));
                                    // main
```

1.4 Pragma Optimization

```
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
1.5 IO Optimization
static inline int gc() {
  constexpr int B = 1 << 20;
  static char buf[B], *p, *q;
  if(p == a \&\&
    (q=(p=buf)+fread(buf,1,B,stdin)) == buf)
   return EOF;
  return *p++;
template < typename T >
static inline bool gn( T &x ) {
 int c = gc(); T sgn = 1; x = 0;
while(('0'>c||c>'9') && c!=EOF && c!='-') c = gc();
if(c == '-') sgn = -1, c = gc();
 if(c == EOF) return false;
 while('0'<=c&&c<='9') x = x*10 + c - '0', c = gc();
 return x *= sgn, true;
```

#pragma GCC optimize("Ofast,no-stack-protector")

#pragma GCC optimize("no-math-errno,unroll-loops")

2 Data Structure

2.1 Dark Magic

2.2 Link-Cut Tree

p->ch[dir]=c;

```
struct Node{
Node *par, *ch[2];
int xor_sum, v;
bool is_rev;
Node(int _v){
 v=xor_sum=_v;is_rev=false;
 par=ch[0]=ch[1]=nullptr;
inline void set_rev(){is_rev^=1;swap(ch[0],ch[1]);}
inline void down(){
 if(is_rev){
  if(ch[0]!=nullptr) ch[0]->set_rev();
   if(ch[1]!=nullptr) ch[1]->set_rev();
   is_rev=false;
 }
inline void up(){
 xor_sum=v;
  if(ch[0]!=nullptr){
  xor_sum^=ch[0]->xor_sum;
  ch[0]->par=this;
 if(ch[1]!=nullptr){
  xor_sum^=ch[1]->xor_sum;
  ch[1]->par=this;
inline bool is_root(){
 return par==nullptr ||\
   (par->ch[0]!=this && par->ch[1]!=this);
bool is_rch(){return !is_root() && par->ch[1]==this;}
} *node[maxn], *stk[maxn];
int top;
void to_child(Node* p,Node* c,bool dir){
```

```
p->up();
inline void rotate(Node* node){
 Node* par=node->par;
 Node* par_par=par->par;
 bool dir=node->is_rch()
 bool par_dir=par->is_rch()
 to_child(par, node->ch[!dir], dir);
 to_child(node,par,!dir);
 if(par_par!=nullptr && par_par->ch[par_dir]==par)
  to_child(par_par,node,par_dir);
 else node->par=par_par;
inline void splay(Node* node){
 Node* tmp=node;
 stk[top++]=node;
 while(!tmp->is_root()){
  tmp=tmp->par;
  stk[top++]=tmp;
 while(top) stk[--top]->down();
 for(Node *fa=node->par;
  !node->is_root();
  rotate(node), fa=node->par)
  if(!fa->is_root())
   rotate(fa->is_rch()==node->is_rch()?fa:node);
inline void access(Node* node){
 Node* last=nullptr;
 while(node!=nullptr){
  splay(node);
  to_child(node, last, true);
  last=node;
  node=node->par;
inline void change_root(Node* node){
 access(node);splay(node);node->set_rev();
inline void link(Node* x, Node* y){
 change_root(x);splay(x);x->par=y;
inline void split(Node* x,Node* y){
 change_root(x);access(y);splay(x);
 to_child(x,nullptr,true);y->par=nullptr;
inline void change_val(Node* node,int v){
access(node);splay(node);node->v=v;node->up();
inline int query(Node* x,Node* y){
 change_root(x);access(y);splay(y);
 return y->xor_sum;
inline Node* find_root(Node* node){
 access(node);splay(node);
 Node* last=nullptr:
 while(node!=nullptr){
  node->down();last=node;node=node->ch[0];
 return last;
set<pii> dic;
inline void add_edge(int u,int v){
 if(u>v) swap(u,v)
 if(find_root(node[u])==find_root(node[v])) return;
 dic.insert(pii(u,v))
link(node[u],node[v]);
inline void del_edge(int u,int v){
 if(u>v) swap(u,v);
 if(dic.find(pii(u,v))==dic.end()) return;
 dic.erase(pii(u,v))
 split(node[u],node[v]);
2.3 LiChao Segment Tree
struct Line{
 int m, k, id;
 Line() : id( -1 ) {}
Line('int a, int'b,'int c')
: m(a), k(b), id(c) {}
```

int at(int x) { return m * x + k; }

```
private:
class LiChao {
                                                              vector< vector< T > > tbl;
private:
                                                              vector< int > lg;
                                                              T cv(Ta, Tb) {
  int n; vector< Line > nodes;
  inline int lc( int x ) { return 2 * x + 1; }
                                                               return Cmp_()( a, b ) ? a : b;
  inline int rc( int x ) { return 2 * x + 2; }
  void insert( int 1, int r, int id, Line ln ) {
                                                             public:
   int m = (1 + r) >> 1;
                                                              void init( T arr[], int n ) {
   if ( nodes[ id ].id == -1 ) {
                                                               // 0-base
   nodes[ id ] = ln;
                                                               lg.resize(n+1);
                                                               lg[0] = -1;
    return:
                                                               for( int i=1 ; i<=n ; ++i ) lg[i] = lg[i>>1] + 1;
   bool atLeft = nodes[ id ].at( 1 ) < ln.at( 1 );</pre>
                                                               tbl.resize(lg[n] + 1);
   if ( nodes[ id ].at( m ) < ln.at( m ) ) {</pre>
                                                               tbl[ 0 ].resize( n );
                                                               copy( arr, arr + n, tbl[ 0 ].begin() );
   atLeft ^= 1; swap( nodes[ id ], ln );
                                                               for ( int i = 1 ; i <= lg[ n ] ; ++ i ) {
  int len = 1 << ( i - 1 ), sz = 1 << i;</pre>
   if ( r - l == 1 ) return;
                                                                tbl[ i ].resize( n - sz + 1 );
   if ( atLeft ) insert( l, m, lc( id ), ln );
                                                                for (int_j = 0; j \le n - sz; ++ j
   else insert( m, r, rc( id ), ln );
                                                                 tbl[i][j] = cv(tbl[i-1][j], tbl[i-1][j+len]);
  int query( int 1, int r, int id, int x ) {
   int ret = 0;
   if ( nodes[ id ].id != -1 )
                                                              T query( int 1, int r ) {
                                                               // 0-base [1, r)
   ret = nodes[ id ].at( x );
   int m = (1 + r) >> 1;
                                                               int wh = lg[ r - l ], len = 1 << wh;</pre>
                                                               return cv( tbl[ wh ][ 1 ], tbl[ wh ][ r - len ] );
   if ( r - l == 1 ) return ret;
   else if (x < m )
                                                            };
    return max( ret, query( 1, m, lc( id ), x ) );
   else
                                                             2.6
                                                                  Linear Basis
    return max( ret, query( m, r, rc( id ), x ) );
                                                             template <int BITS>
public:
                                                             struct LinearBasis {
 void build( int n_ ) {
                                                              array<uint64_t, BITS> basis;
                                                              Basis() { basis.fill(0); }
  n = n_; nodes.clear();
  nodes.resize( n << 2, Line() );</pre>
                                                              void add(uint64_t x)
                                                               for (int i = 0; i < BITS; ++i) if ((x >> i) & 1) {
 void insert( Line ln ) { insert( 0, n, 0, ln ); }
                                                                if (basis[i] == 0) {
  int query( int x ) { return query( 0, n, 0, x ); }
                                                                 basis[i] = x;
} lichao;
                                                                 return;
2.4 Treap
                                                                x ^= basis[i];
namespace Treap{
                                                               }
#define sz( x ) ( ( x ) ? ( ( x )->size ) : 0 )
                                                              bool ok(uint64_t x) {
struct node{
                                                               for (int i = 0; i < BITS; ++i)</pre>
 int size;
  uint32_t pri;
                                                                if ((x >> i) & 1) x ^= basis[i];
 node *lc, *rc;
node() : size(0), pri(rand()), lc(0), rc(0) {}
                                                               return x == 0;
                                                            }:
  void pull() {
  size = 1;
                                                                   Binary Search On Segment Tree
  if ( lc ) size += lc->size;
   if ( rc ) size += rc->size;
                                                            // find_first = x -> minimal x s.t. check( [a, x) )
                                                             // find_last = x \rightarrow maximal x s.t. check([x, b))
  }
                                                             template <typename C>
node* merge( node* L, node* R ) {
  if ( not L or not R ) return L ? L : R;
                                                             int find_first(int 1, const C &check) {
                                                              if (1 >= n)
 if ( L->pri > R->pri ) {
                                                               return n;
  L->rc = merge( L->rc, R ); L->pull();
                                                              1 += sz;
   return L;
                                                              for (int i = height; i > 0; i--)
 } else {
                                                               propagate(1 >> i);
  R->lc = merge( L, R->lc ); R->pull();
                                                              Monoid sum = identity;
                                                              do {
   return R;
                                                               while ((1 \& 1) == 0)
}
                                                                1 >>= 1
void split_by_size( node*rt,int k,node*&L,node*&R ) {
                                                               if (check(f(sum, data[1]))) {
 if ( not rt ) L = R = nullptr;
                                                                while (1 < sz) {</pre>
  else if( sz( rt->lc ) + 1 <= k ) {
                                                                 propagate(1);
  split_by_size( rt->rc,k-sz(rt->lc)-1,L->rc,R );
                                                                 auto nxt = f(sum, data[1]);
  L->pull();
                                                                 if (not check(nxt)) {
 } else {
                                                                  sum = nxt;
  R = rt:
                                                                  1++;
   split_by_size( rt->lc, k, L, R->lc );
   R->pull();
 }
                                                                return 1 + 1 - sz;
#undef sz
                                                               sum = f(sum, data[1++]);
                                                              } while ((1 & -1) != 1);
                                                              return n;
2.5 Sparse Table
template < typename T, typename Cmp_ = less< T > >
class SparseTable {
                                                            template <typename C>
```

if (not ins[t]) {

st.push_back(t);

```
int find_last(int r, const C &check) {
                                                                  ins[t] = true;
if (r <= 0)
                                                                 if (dfn[v]) {
 return -1;
                                                                  low[u] = min(low[u], dfn[v]);
 r += sz;
for (int i = height; i > 0; i--)
                                                                  continue
 propagate((r - 1) >> i);
                                                                 } ++ch; dfs(v, u)
Monoid sum = identity;
                                                                 low[u] = min(low[u], low[v]);
                                                                 if (low[v] >= dfn[u]) {
do {
                                                                  ap[u] = true;
 while (r > 1 \text{ and } (r \& 1))
                                                                  while (true) {
                                                                   int eid = st.back(); st.pop_back();
  r >>= 1
  if (check(f(data[r], sum))) {
                                                                   bcc[eid] = ecnt;
                                                                   if (eid == t) break;
  while (r < sz) {</pre>
   propagate(r);
    r = (r << 1) + 1;
                                                                  ecnt++;
   auto nxt = f(data[r], sum);
    if (not check(nxt)) {
    sum = nxt;
                                                                if (ch == 1 and u == f) ap[u] = false;
                                                               }
     r--;
   }
                                                              public:
                                                               void init(int n_) {
   return r - sz;
                                                                G.clear(); G.resize(n = n_);
                                                                ecnt = 0; ap.assign(n, false);
 sum = f(data[r], sum);
                                                                low.assign(n, 0); dfn.assign(n, 0);
} while ((r & -r) != r);
return -1;
                                                               void add_edge(int u, int v) {
                                                               G[u].emplace_back(v, ecnt);
                                                                G[v].emplace_back(u, ecnt++);
3
    Graph
                                                               void solve() {
                                                                ins.assign(ecnt, false);
   BCC Edge
                                                                bcc.resize(ecnt); ecnt = 0;
class BCC_Bridge {
                                                                for (int i = 0; i < n; ++i)</pre>
private:
                                                                 if (not dfn[i]) dfs(i, i);
 int n. ecnt:
 vector<vector<pair<int,int>>> G;
                                                               int get_id(int x) { return bcc[x]; }
  vector<int> dfn, low;
                                                               int count() { return ecnt; }
 vector<bool> bridge;
                                                               bool is_ap(int x) { return ap[x]; }
  void dfs(int u, int f)
                                                            } bcc_ap;
  dfn[u] = low[u] = dfn[f] + 1;
  for (auto [v, t]: G[u]) {
                                                             3.3 2-SAT (SCC)
    if (v == f) continue;
                                                            class TwoSat{
    if (dfn[v]) {
                                                              private:
    low[u] = min(low[u], dfn[v]);
                                                               int n:
     continue;
                                                               vector<vector<int>> rG,G,sccs;
                                                               vector<int> ord,idx;
   dfs(v, u);
                                                               vector<bool> vis,result;
   low[u] = min(low[u], low[v]);
                                                               void dfs(int u){
    if (low[v] > dfn[u]) bridge[t] = true;
                                                                vis[u]=true
                                                                for(int v:G[u])
                                                                 if(!vis[v]) dfs(v);
public:
                                                                ord.push_back(u);
  void init(int n_) {
  G.clear(); G.resize(n = n_);
                                                               void rdfs(int u){
  low.assign(n, ecnt = 0);
                                                                vis[u]=false;idx[u]=sccs.size()-1;
  dfn.assign(n, 0);
                                                                sccs.back().push_back(u);
                                                                for(int v:rG[u])
  void add_edge(int u, int v) {
                                                                 if(vis[v])rdfs(v);
  G[u].emplace_back(v, ecnt);
G[v].emplace_back(u, ecnt++);
                                                              public:
                                                               void init(int n_){
 void solve() {
                                                                n=n_;G.clear();G.resize(n);
  bridge.assign(ecnt, false);
                                                                rG.clear();rG.resize(n);
  for (int i = 0; i < n; ++i)</pre>
                                                                sccs.clear();ord.clear();
   if (not dfn[i]) dfs(i, i);
                                                                idx.resize(n);result.resize(n);
 bool is_bridge(int x) { return bridge[x]; }
                                                               void add_edge(int u,int v){
} bcc_bridge;
                                                                G[u].push_back(v);rG[v].push_back(u);
3.2 BCC Vertex
                                                               void orr(int x,int y){
class BCC_AP {
                                                                if ((x^y)==1) return
                                                                add_edge(x^1,y); add_edge(y^1,x);
private:
 int n, ecnt;
                                                               bool solve(){
 vector<vector<pair<int,int>>> G;
 vector<int> bcc, dfn, low, st;
                                                                vis.clear();vis.resize(n);
 vector<bool> ap, ins;
                                                                for(int i=0;i<n;++i)</pre>
 void dfs(int u, int f) {
  dfn[u] = low[u] = dfn[f] + 1;
                                                                 if(not vis[i])dfs(i);
                                                                reverse(ord.begin(),ord.end());
  int ch = 0;
                                                                for (int u:ord){
   for (auto [v, t]: G[u]) if (v != f) {
                                                                 if(!vis[u])continue;
```

sccs.push_back(vector<int>());

rdfs(u);

time_ = 0; dfschain(1, 1);

```
for(int i=0;i<n;i+=2)</pre>
                                                                PII get_subtree(int u) { return {tl[ u ],tr[ u ] }; }
    if(idx[i]==idx[i+1])
                                                                vector< PII > get_path( int u , int v ){
                                                                 vector< PII > res;
     return false
                                                                 int g = lca( u, v );
   vector<bool> c(sccs.size());
   for(size_t i=0;i<sccs.size();++i){</pre>
                                                                 while ( chain[ u ] != chain[ g ] ) {
                                                                  int s = chain_st[ chain[ u ] ];
res.emplace_back( tl[ s ], tl[ u ] + 1 );
    for(size_t j=0;j<sccs[i].size();++j){</pre>
     result[sccs[i][j]]=c[i];
     c[idx[sccs[i][j]^1]]=!c[i];
                                                                  u = fa[ s ][ 0 ];
                                                                 res.emplace\_back( tl[ g ], tl[ u ] + 1 );
                                                                 while ( chain[ v ] != chain[ g ] ) {
   return true;
                                                                  int s = chain_st[ chain[ v ] ]
 bool get(int x){return result[x];}
                                                                  res.emplace_back( tl[ s ], tl[ v ] + 1 );
  inline int get_id(int x){return idx[x];}
                                                                  v = fa[ s ][ 0 ];
  inline int count(){return sccs.size();}
                                                                 res.emplace_back( tl[ g ] + 1, tl[ v ] + 1 );
                                                                 return res:
3.4 Lowbit Decomposition
                                                                 /* res : list of intervals from u to v
class LowbitDecomp{
                                                                  * ( note only nodes work, not edge )
                                                                  * usage :
private:
                                                                  * vector< PII >& path = tree.get_path( u , v )
int time_, chain_, LOG_N;
                                                                  * for( auto [ 1, r ] : path ) {
vector< vector< int > > G, fa;
vector< int > tl, tr, chain, chain_st;
                                                                  * 0-base [ 1, r )
                                                                  * }
// chain_ : number of chain
                                                                  */
// tl, tr[ u ] : subtree interval in the seq. of u
                                                                }
 // chain_st[ u ] : head of the chain contains u
// chian[ u ] : chain id of the chain u is on
                                                              } tree;
void predfs( int u, int f ) {
                                                               3.5
                                                                     MaxClique
  chain[u] = 0;
  for ( int v : G[ u ] ) {
                                                               // contain a self loop u to u, than u won't in clique
  if ( v == f ) continue;
                                                               template < size_t MAXN >
   predfs( v, u );
                                                               class MaxClique{
   if( lowbit( chain[ u ] ) < lowbit( chain[ v ] ) )</pre>
                                                               private:
    chain[ u ] = chain[ v ];
                                                                using bits = bitset< MAXN >;
                                                                bits popped, G[ MAXN ], ans;
size_t deg[ MAXN ], deo[ MAXN ], n;
 if ( not chain[ u ] )
   chain[ u ] = chain_ ++;
                                                                void sort_by_degree() {
                                                                 popped.reset();
 void dfschain( int u, int f ) {
                                                                 for ( size_t i = 0 ; i < n ; ++ i )</pre>
 fa[ u ][ 0 ] = f;
for ( int i = 1 ; i < LOG_N ; ++ i )
                                                                   deg[ i ] = G[ i ].count();
                                                                 for ( size_t i = 0 ; i < n ; ++ i ) {
  fa[u][i] = fa[fa[u][i-1]][i-1];
                                                                   size_t mi = MAXN, id = 0;
                                                                   for ( size_t j = 0 ; j < n ; ++ j )
  if ( not popped[ j ] and deg[ j ] < mi )
    mi = deg[ id = j ];</pre>
  tl[ u ] = time_++;
  if ( not chain_st[ chain[ u ] ] )
   chain_st[ chain[ u ] ] = u;
  for ( int v : G[ u ] )
                                                                   popped[ deo[ i ] = id ] = 1;
  if ( v != f and chain[ v ] == chain[ u ] )
                                                                   for( size_t u = G[ i ]._Find_first() ;
                                                                    u < n ; u = G[ i ]._Find_next( u ) )</pre>
    dfschain( v, u );
  for ( int v : G[ u ] )
                                                                     -- deg[ u ];
   if ( v != f and chain[ v ] != chain[ u ] )
    dfschain( v, u );
                                                                void BK( bits R, bits P, bits X ) {
  tr[ u ] = time_;
                                                                 if (R.count()+P.count() <= ans.count()) return;</pre>
bool anc( int u, int v ) {
  return tl[ u ] <= tl[ v ] and tr[ v ] <= tr[ u ];</pre>
                                                                 if ( not P.count() and not X.count() ) {
                                                                  if ( R.count() > ans.count() ) ans = R;
                                                                  return;
public:
                                                                 }
int lca( int u, int v ) {
                                                                 /* greedily chosse max degree as pivot
  if ( anc( u, v ) ) return u;
                                                                 bits cur = P | X; size_t pivot = 0, sz = 0;
 for ( int i = LOG_N - 1 ; i >= 0 ; -- i )
if ( not anc( fa[ u ][ i ], v ) )
                                                                 for ( size_t u = cur._Find_first() ;
                                                                  u < n ; u = cur._Find_next( u )</pre>
   u = fa[ u ][ i ];
                                                                   if ( deg[ u ] > sz ) sz = deg[ pivot = u ];
                                                                 cur = P & ( ~G[ pivot ] );
  return fa[ u ][ 0 ];
                                                                 */ // or simply choose first
void init( int n ) {
                                                                 bits cur = P & (~G[ ( P | X )._Find_first() ]);
 fa.assign( ++n, vector< int >( LOG_N ) );
                                                                 for ( size_t u = cur._Find_first()
                                                                  u < n ; u = cur._Find_next( u ) ) {
if ( R[ u ] ) continue;</pre>
  for ( LOG_N = 0 ; ( 1 << LOG_N ) < n ; ++ LOG_N );
 G.clear(); G.resize( n );
 tl.assign( n, 0 ); tr.assign( n, 0 );
                                                                  R[u] = 1;
 chain.assig( n, 0 ); chain_st.assign( n, 0 );
                                                                  BK( R, P & G[ u ], X & G[ u ] );
                                                                  R[u] = P[u] = 0, X[u] = 1;
 void add_edge( int u , int v ) {
  // 1-base
 G[ u ].push_back( v );
                                                               public:
 G[ v ].push_back( u );
                                                                void init( size_t n_ ) {
                                                                 n = n_{-};
void decompose(){
                                                                 for ( size_t i = 0 ; i < n ; ++ i )
 chain_ = 1;
                                                                  G[ i ].reset();
 predfs( 1, 1 );
                                                                 ans.reset();
```

void add_edges(int u, bits S) { G[u] = S; }

```
void add_edge( int u, int v ) {
                                                                sort(r.begin(), r.end(),
                                                                 [&](int i, int j) { return d[i] > d[j]; });
  G[u][v] = G[v][u] = 1;
                                                                csort(r, c);
                                                                dfs(r, c, 1, mask);
return ans; // sol[0 ~ ans-1]
 int solve() {
  sort_by_degree(); // or simply iota( deo... )
  for ( size_t i = 0 ; i < n ; ++ i )</pre>
                                                               }
   deg[ i ] = G[ i ].count();
                                                              } graph;
  bits pob, nob = 0; pob.set();
                                                              3.7 Virtural Tree
  for (size_t i=n; i<MAXN; ++i) pob[i] = 0;</pre>
  for ( size_t i = 0 ; i < n ; ++ i ) {
    size_t v = deo[ i ];</pre>
                                                              inline bool cmp(const int &i, const int &j) {
                                                               return dfn[i] < dfn[j];</pre>
   bits tmp; tmp[ v ] = 1;
   BK( tmp, pob & G[ v ], nob & G[ v ] );
pob[ v ] = 0, nob[ v ] = 1;
                                                              void build(int vectrices[], int k) {
                                                               static int stk[MAX_N];
                                                               sort(vectrices, vectrices + k, cmp);
  return static_cast< int >( ans.count() );
                                                               stk[sz++] = 0;
                                                               for (int i = 0; i < k; ++i) {
  int u = vectrices[i], lca = LCA(u, stk[sz - 1]);</pre>
};
                                                                if (lca == stk[sz - 1]) stk[sz++] = u;
3.6 MaxCliqueDyn
                                                                else {
constexpr int kN = 150;
                                                                 while (sz \ge 2 \&\& dep[stk[sz - 2]] \ge dep[lca]) {
struct MaxClique { // Maximum Clique
                                                                  addEdge(stk[sz - 2], stk[sz - 1]);
sz--:
                                                                 if (stk[sz - 1] != lca) {
 void init(int _n) {
 n = n, ans q = 0;
                                                                  addEdge(lca, stk[--sz]);
  for (int i = 0; i < n; i++) a[i].reset();</pre>
                                                                  stk[sz++] = lca, vectrices[cnt++] = lca;
 void addEdge(int u, int v) { a[u][v] = a[v][u] = 1; }
void csort(vector<int> &r, vector<int> &c) {
                                                                 stk[sz++] = u;
 int mx = 1, km = max(ans - q + 1, 1), t = 0,
                                                               for (int i = 0; i < sz - 1; ++i)
    m = int(r.size())
  cs[1].reset(); cs[2].reset()
                                                                addEdge(stk[i], stk[i + 1]);
  for (int i = 0; i < m; i++) {
   int p = r[i], k = 1
                                                              3.8 Centroid Decomposition
   while ((cs[k] & a[p]).count()) k++;
   if (k > mx) cs[++mx + 1].reset();
                                                              struct Centroid {
   cs[k][p] = 1;
                                                               vector<vector<int64_t>> Dist;
                                                               vector<int> Parent, Depth;
   if (k < km) r[t++] = p;
                                                               vector<int64_t> Sub, Sub2;
  c.resize(m);
                                                               vector<int> Sz, Sz2;
  if(t) c[t-1] = 0;
                                                               Centroid(vector<vector<pair<int, int>>> g) {
  for (int k = km; k <= mx; k++) {</pre>
                                                                int N = g.size()
  for (int p = int(cs[k]._Find_first());
                                                                vector<bool> Vis(N);
      p < kN; p = int(cs[k]._Find_next(p))) {
                                                                vector<int> sz(N), mx(N);
    r[t] = p; c[t++] = k;
                                                                vector<int> Path;
                                                                Dist.resize(N)
  }
                                                                Parent.resize(N);
                                                                Depth.resize(N)
 void dfs(vector<int> &r, vector<int> &c, int 1,
                                                                auto DfsSz = [\&](auto dfs, int x) -> void {
  bitset<kN> mask) {
                                                                 Vis[x] = true; sz[x] = 1; mx[x] = 0;
                                                                 for (auto [u, w] : g[x]) {
  while (!r.empty()) {
                                                                  if (Vis[u]) continue;
   int p = r.back(); r.pop_back();
                                                                  dfs(dfs, u)
   mask[p] = 0;
   if (q + c.back() <= ans) return;</pre>
                                                                  sz[x] += sz[u];
                                                                  mx[x] = max(mx[x], sz[u]);
   cur[q++] = p;
   vector<int> nr, nc;
   bitset<kN> nmask = mask & a[p];
                                                                 Path.push_back(x);
   for (int i : r)
                                                                };
    if (a[p][i]) nr.push_back(i);
                                                                auto DfsDist = [&](auto dfs, int x, int64_t D = 0)
                                                                 -> void {
   if (!nr.empty()) {
    if (1 < 4) {
                                                                 Dist[x].push_back(D); Vis[x] = true;
     for (int i : nr)
                                                                 for (auto [u, w] : g[x]) {
      d[i] = int((a[i] & nmask).count());
                                                                  if (Vis[u]) continue;
     sort(nr.begin(), nr.end(),
                                                                  dfs(dfs, u, D + w);
      [&](int x, int y)
                                                                 }
       return d[x] > d[y];
                                                                };
                                                                auto Dfs = [&]
      });
                                                                 (auto dfs, int x, int D = 0, int p = -1)->void {
   csort(nr, nc); dfs(nr, nc, 1 + 1, nmask);
} else if (q > ans) {
                                                                 Path.clear(); DfsSz(DfsSz, x);
                                                                 int M = Path.size();
                                                                 int C = -1;
    ans = q; copy(cur, cur + q, sol);
                                                                 for (int u : Path) {
   c.pop_back(); q--;
                                                                  if (max(M - sz[u], mx[u]) * 2 <= M) C = u;
  }
                                                                  Vis[u] = false;
 int solve(bitset<kN> mask) { // vertex mask
                                                                 DfsDist(DfsDist, C);
                                                                 for (int u : Path) Vis[u] = false;
  vector<int> r, c;
  for (int i = 0; i < n; i++)
                                                                 Parent[C] = p; Vis[C] = true;
  if (mask[i]) r.push_back(i);
for (int i = 0; i < n; i++)</pre>
                                                                 Depth[C] = D;
                                                                 for (auto [u, w] : g[C]) {
                                                                  if (Vis[u]) continue;
   d[i] = int((a[i] & mask).count());
```

```
dfs(dfs, u, D + 1, C);
                                                                  FZ(vst);edgeID.clear();cycle.clear();rho.clear();
   }
                                                                  for (int i=n; !vst[st]; st=prv[i--][st]) {
  Dfs(Dfs, 0); Sub.resize(N); Sub2.resize(N);
                                                                   vst[st]++
                                                                   edgeID.PB(prve[i][st]);
  Sz.resize(N); Sz2.resize(N);
                                                                   rho.PB(st);
 void Mark(int v) {
  int x = v, z = -1;
                                                                  while (vst[st] != 2) {
  for (int i = Depth[v]; i >= 0; --i) {
                                                                   int v = rho.back(); rho.pop_back();
                                                                   cycle.PB(v);
   Sub[x] += Dist[v][i]; Sz[x]++;
   if (z != -1) {
                                                                   vst[v]++;
    Sub2[z] += Dist[v][i];
    Sz2[z]++;
                                                                  reverse(ALL(edgeID));
                                                                  edgeID.resize(SZ(cycle));
   z = x; x = Parent[x];
                                                                  return mmc;
  }
                                                               } mmc;
 int64_t Query(int v) {
                                                                3.11 Mo's Algorithm on Tree
 int64_t res = 0;
  int x = v, z = -1
                                                                int q; vector< int > G[N];
 for (int i = Depth[v]; i >= 0; --i) {
  res += Sub[x] + 1LL * Sz[x] * Dist[v][i];
                                                                struct Que{
                                                                int u, v, id;
} que[ N ];
   if (z != -1) res-=Sub2[z]+1LL*Sz2[z]*Dist[v][i];
                                                                int dfn[N], dfn_, block_id[N], block_, stk[N], stk_;
   z = x; x = Parent[x];
                                                                void_dfs( int u, int f ) {
                                                                 dfn[ u ] = dfn_++; int saved_rbp = stk_;
for ( int v : G[ u ] ) {
  return res;
                                                                  if ( v == f ) continue;
};
                                                                  dfs( v, u );
3.9
     Tree Hashing
                                                                  if ( stk_ - saved_rbp < SQRT_N ) continue;</pre>
                                                                  for ( ++ block_ ; stk_ != saved_rbp ; )
  block_id[ stk[ -- stk_ ] ] = block_;
uint64_t hsah(int u, int f) {
 uint64_t r = 127;
 for (int v : G[ u ]) if (v != f) {
 uint64_t hh = hsah(v, u);
                                                                stk[ stk_ ++ ] = u;
  r=(r+(hh*hh)%1010101333)%1011820613;
                                                                bool inPath[ N ];
                                                                void Diff( int u ) {
return r;
                                                                if ( inPath[ u ] ^= 1 ) { /*remove this edge*/ }
                                                                 else { /*add this edge*/ }
3.10 Minimum Mean Cycle
/* minimum mean cycle O(VE) */
                                                                void traverse( int& origin_u, int u ) {
                                                                for ( int g = lca( origin_u, u )
struct MMC{
                                                                  origin_u != g ; origin_u = parent_of[ origin_u ] )
#define FZ(n) memset((n),0,sizeof(n))
#define E 101010
                                                                   Diff( origin_u );
#define V 1021
                                                                 for (int v = u; v != origin_u; v = parent_of[v])
#define inf 1e9
                                                                  Diff( v );
 struct Edge { int v,u; double c; };
                                                                 origin_u = u;
 int n, m, prv[V][V], prve[V][V], vst[V];
                                                                }
 Edge e[E];
                                                                void solve() {
 vector<int> edgeID, cycle, rho;
                                                                 dfs( 1, 1 );
 double d[V][V];
                                                                 while ( stk_ ) block_id[ stk[ -- stk_ ] ] = block_;
 void init( int _n ) { n = _n; m = 0; }
// WARNING: TYPE matters
                                                                 sort( que, que + q, [](const Que& x, const Que& y) {
                                                                  return tie( block_id[ x.u ], dfn[ x.v ] )
 void add_edge( int vi , int ui , double ci )
                                                                      < tie( block_id[ y.u ], dfn[ y.v ] );
 { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
                                                                 } );
                                                                 int U = 1, V = 1;
                                                                 for ( int i = 0 ; i < q ; ++ i ) {
  for(int i=0; i<n; i++) d[0][i]=0;
  for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);</pre>
                                                                  pass( U, que[ i ].u );
pass( V, que[ i ].v );
   for(int j=0; j<m; j++) {</pre>
                                                                  // we could get our answer of que[ i ].id
    int v = e[j].v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
                                                               }
                                                                /*
     d[i+1][u] = d[i][v]+e[j].c;
     prv[i+1][u] = v;
                                                                Method 2:
     prve[i+1][u] = j;
                                                               dfs u:
                                                                push u
                                                                 iterate subtree
                                                                Let P = LCA(u, v), and St(u) \le St(v)
                                                               if (P == u) query[St(u), St(v)]
 double solve(){
  // returns inf if no cycle, mmc otherwise
                                                                else query[Ed(u), St(v)], query[St(P), St(P)]
  double mmc=inf;
  int st = -1
                                                                3.12 Minimum Steiner Tree
  bellman_ford();
  for(int i=0; i<n; i++) {</pre>
                                                               // Minimum Steiner Tree
   double avg=-inf;
                                                               // 0(V 3^T + V^2 2^T)
                                                               struct SteinerTree{
   for(int k=0; k<n; k++) {</pre>
    if(d[n][i]<inf-eps)</pre>
                                                               #define V 33
     avg=max(avg,(d[n][i]-d[k][i])/(n-k));
                                                                #define T 8
                                                               #define INF 1023456789
    else avg=max(avg,inf);
                                                                 int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
                                                                void init( int _n ){
   if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
```

if (in[i] == -inf)
return -inf;

// find cycle

```
int tot = 0;
  for( int i = 0 ; i < n ; i ++ ){</pre>
                                                                  vector<int> id(n, -1), vis(n, -1);
   for( int j = 0 ; j < n ; j ++ )</pre>
                                                                   for (int i = 0; i < n; i++) {
   dst[ i ][ j ] = INF;
dst[ i ][ i ] = 0;
                                                                    ans += in[i];
                                                                    for (int x = i; x != -1 && id[x] == -1; x = prv[x])
  }
                                                                     if (vis[x] == i) {
                                                                     for (int y = prv[x]; y != x; y = prv[y])
 void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
                                                                       id[y] = tot;
  dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
                                                                      id[x] = tot++;
                                                                      break:
 void shortest_path(){
  for( int k = 0 ; k < n ; k ++ )</pre>
                                                                     vis[x] = i;
   for( int i = 0 ; i < n ; i ++ )</pre>
                                                                   }
    for( int j = 0 ; j < n ; j ++ )</pre>
     dst[ i ][ j ] = min( dst[ i ][ j ],
    dst[ i ][ k ] + dst[ k ][ j ] );
                                                                  if (!tot)
                                                                    return ans;
                                                                   for (int i = 0; i < n; i++)
 int solve( const vector<int>& ter ){
                                                                   if (id[i] == -1)
  int t = (int)ter.size();
                                                                     id[i] = tot++;
  for( int i = 0 ; i < ( 1 << t ) ; i ++ )
                                                                   // shrink
   for( int j = 0; j < n; j ++ )
dp[ i ][ j ] = INF;
                                                                  for (auto &e : E) {
  if (id[e.u] != id[e.v])
  for( int i = 0 ; i < n ; i ++ )
                                                                    e.w -= in[e.v];
   dp[0][i] = 0;
                                                                   e.u = id[e.u], e.v = id[e.v];
  for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){
  if( msk == ( msk & (-msk) ) ){</pre>
                                                                  n = tot:
    int who = __lg( msk );
                                                                  root = id[root];
    for( int i = 0 ; i < n ; i ++ )
dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
                                                                 assert(false);
    continue:
                                                               } DMST;
   for( int i = 0 ; i < n ; i ++ )</pre>
                                                               3.14
                                                                      Dominator Tree
    for( int submsk = ( msk - 1 ) & msk ; submsk ;
         submsk = ( submsk - 1 ) & msk )
                                                               namespace dominator {
      vector<int> g[maxn], r[maxn], rdom[maxn];
                                                               int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn];
               dp[ msk ^ submsk ][ i ] );
                                                               int dom[maxn], val[maxn], rp[maxn], tk;
   for( int i = 0 ; i < n ; i ++ ){</pre>
                                                               void init(int n) {
    tdst[ i ] = INF;
                                                                // vertices are numbered from 0 to n - 1
    fill(dfn, dfn + n, -1);fill(rev, rev + n, -1);
                                                                fill(fa, fa + n, -1); fill(val, val + n, -1);
                                                                fill(sdom, sdom + n, -1); fill(rp, rp + n, -1);
                                                                fill(dom, dom + n, -1); tk = \hat{0};
   for( int i = 0 ; i < n ; i ++ )</pre>
                                                                for (int i = 0; i < n; ++i) {
    dp[ msk ][ i ] = tdst[ i ];
                                                                 g[i].clear(); r[i].clear(); rdom[i].clear();
  int ans = INF:
  for( int i = 0 ; i < n ; i ++ )</pre>
                                                               void add_edge(int x, int y) { g[x].push_back(y); }
   ans = min( ans , dp[ (1 << t) - 1 ][ i ] );
                                                               void dfs(int x) {
                                                                rev[dfn[x] = tk] = x;
  return ans;
                                                                fa[tk] = sdom[tk] = val[tk] = tk; tk ++;
} solver;
                                                                for (int u : g[x]) {
                                                                 if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
      Directed Minimum Spanning Tree
                                                                 r[dfn[u]].push_back(dfn[x]);
struct DirectedMST { // find maximum
 struct Edge {
  int u, v;
                                                               void merge(int x, int y) { fa[x] = y; }
                                                               int find(int x, int c = 0) {
  int w:
                                                                if (fa[x] == x) return c ? -1 : x;
  Edge(int u, int v, int w) : u(u), v(v), w(w) {}
                                                                int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
 vector<Edge> Edges;
 void clear() { Edges.clear(); }
                                                                if (sdom[val[x]]>sdom[val[fa[x]]]) val[x]=val[fa[x]];
 void addEdge(int a, int b, int w) { Edges.emplace_back
                                                                fa[x] = p;
return c ? p : val[x];
    (a, b, w); }
 int solve(int root, int n) {
                                                               vector<int> build(int s, int n) {
// return the father of each node in the dominator tree
  vector<Edge> E = Edges;
  int ans = 0;
                                                               // p[i] = -2 if i is unreachable from s
  while (true) {
   // find best in edge
                                                                dfs(s);
   vector<int> in(n, -inf), prv(n, -1);
                                                                for (int i = tk - 1; i >= 0; --i) {
  for (int u:r[i]) sdom[i]=min(sdom[i],sdom[find(u)]);
   for (auto e : E)
    if (e.u != e.v && e.w > in[e.v]) {
                                                                 if (i) rdom[sdom[i]].push_back(i);
                                                                 for (int &u : rdom[i]) {
     in[e.v] = e.w;
                                                                  int p = find(u);
     prv[e.v] = e.u;
                                                                  if (sdom[p] == i) dom[u] = i;
else dom[u] = p;
   in[root] = 0;
   prv[root] = -1;
   for (int i = 0; i < n; i++)
                                                                 if (i) merge(i, rp[i]);
```

vector<int> p(n, -2); p[s] = -1;

for (int i = 1; i < tk; ++i)

```
if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
 for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
 return p;
3.15
      Edge Coloring
// max(d_u) + 1 edge coloring, time: O(NM)
int C[kN][kN], G[kN][kN]; // 1-based, G: ans
void clear(int N) {
 for (int i = 0; i <= N; i++)
  for (int j = 0; j <= N; j++)</pre>
    C[i][j] = G[i][j] = 0;
void solve(vector<pair<int, int>> &E, int N) {
int X[kN] = {}, a;
auto update = [&](int u) {
  for (X[u] = 1; C[u][X[u]]; X[u]++);
 auto color = [&](int u, int v, int c) {
  int p = G[u][v];
  G[u][v] = G[v][u] = c;
  C[u][c] = v, C[v][c] = u;
  C[u][p] = C[v][p] = 0;
  if (p) X[u] = X[v] = p;
  else update(u), update(v);
  return p;
 } ;
 auto flip = [&](int u, int c1, int c2) {
  int p = C[u][c1];
  swap(C[u][c1], C[u][c2]);
 if (p) G[u][p] = G[p][u] = c2;
if (!C[u][c1]) X[u] = c1;
  if (!C[u][c2]) X[u] = c2;
  return p;
 };
 for (int i = 1; i <= N; i++) X[i] = 1;
 for (int t = 0; t < E.size(); t++) {</pre>
  auto [u, v] = E[t];
  int v0 = v, c = X[u], c0 = c, d;
  vector<pair<int, int>> L; int vst[kN] = {};
  while (!G[u][v0]) {
   L.emplace_back(v, d = X[v]);
   if (!C[v][c]) for(a=L.size()-1;a>=0;a--)
   c = color(u, L[a].first, c);
else if(!C[u][d])for(a=L.size()-1;a>=0;a--)
     color(u, L[a].first, L[a].second);
   else if (vst[d]) break
   else vst[d] = 1, v = C[u][d];
  if (!G[u][v0]) {
   for (; v; v = flip(v, c, d), swap(c, d));
   if (C[u][c0]) { a = int(L.size()) - 1;
    while (--a >= 0 && L[a].second != c);
    for(;a>=0;a--)color(u,L[a].first,L[a].second);
   } else t--;
     Matching & Flow
4
    Kuhn Munkres
class KM {
private:
 static constexpr lld INF = 1LL << 60;</pre>
 vector<lld> h1,hr,slk;
```

```
private:
    static constexpr lld INF = 1LL << 60;
    vector<lld> hl,hr,slk;
    vector<int> fl,fr,pre,qu;
    vector<vector<lld> w;
    vector<bool> vl,vr;
    int n, ql, qr;
    bool check(int x) {
        if (vl[x] = true, fl[x] != -1)
            return vr[qu[qr++] = fl[x]] = true;
        while (x != -1) swap(x, fr[fl[x] = pre[x]]);
        return false;
    }
    void bfs(int s) {
        fill(slk.begin(), slk.end(), INF);
        fill(vl.begin(), vl.end(), false);
        fill(vr.begin(), vr.end(), false);
        ql = qr = 0;
```

```
while (true) {
   11d d;
   while (ql < qr) {</pre>
    for (int x = 0, y = qu[ql++]; x < n; ++x) {
     if(!v1[x]&&slk[x]>=(d=h1[x]+hr[y]-w[x][y])){
      if (pre[x] = y, d) slk[x] = d;
       else if (!check(x)) return;
    }
   d = INF;
   for (int x = 0; x < n; ++x)
    if (!v1[x] \&\& d > s1k[x]) d = s1k[x];
   for (int x = 0; x < n; ++x) {
    if (v1[x]) h1[x] += d;
    else slk[x] -= d;
    if (vr[x]) hr[x] -= d;
   for (int x = 0; x < n; ++x)
    if (!v1[x] && !slk[x] && !check(x)) return;
public:
 void init( int n_ ) {
  n = n_; qu.resize(n);
  fl.clear(); fl.resize(n, -1);
fr.clear(); fr.resize(n, -1);
hr.clear(); hr.resize(n); hl.resize(n);
  w.clear(); w.resize(n, vector<lld>(n));
  slk.resize(n); pre.resize(n);
  vl.resize(n); vr.resize(n);
 void set_edge( int u, int v, lld x ) {w[u][v] = x;}
 1ld solve() {
  for (int i = 0; i < n; ++i)</pre>
   hl[i] = *max_element(w[i].begin(), w[i].end());
  for (int i = 0; i < n; ++i) bfs(i);</pre>
  11d res = 0;
  for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
  return res;
 }
} km;
4.2 Bipartite Matching
class BipartiteMatching{
private:
 vector<int> X[N], Y[N];
 int fX[N], fY[N], n;
 bitset<N> walked;
 bool dfs(int x)
  for(auto i:X[x]){
   if(walked[i])continue;
   walked[i]=1;
   if(fY[i]==-1||dfs(fY[i])){
    fY[i]=x;fX[x]=i;
    return 1;
   }
  return 0;
public:
 void init(int _n){
  n=_n; walked.reset();
  for(int i=0;i<n;i++){</pre>
   X[i].clear();Y[i].clear();
   fX[i]=fY[i]=-1;
  }
 void add_edge(int x, int y){
  X[x].push_back(y); Y[y].push_back(y);
 int solve(){
  int cnt = 0;
  for(int i=0;i<n;i++){</pre>
   walked.reset();
   if(dfs(i)) cnt++;
  // return how many pair matched
  return cnt;
```

qu[qr++] = s;

vr[s] = true;

for (int i=0; i<n; i++)</pre>

```
| };
                                                                for (int j=0; j<n; j++)</pre>
                                                                 edge[i][j] = 0;
4.3 General Graph Matching
namespace matching {
                                                              void set_edge(int u, int v, int w) {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
                                                               edge[u][v] = edge[v][u] = w;
vector<int> g[kN];
queue<int> q;
                                                              bool SPFA(int u){
                                                               if (onstk[u]) return true;
void Init(int n) {
                                                               stk.PB(u);
 for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;</pre>
 for (int i = 0; i < n; ++i) g[i].clear();</pre>
                                                               onstk[u] = 1;
                                                               for (int v=0; v<n; v++){</pre>
                                                                if (u != v && match[u] != v && !onstk[v]){
void AddEdge(int u, int v) {
                                                                 int m = match[v]
 g[u].push_back(v);
 g[v].push_back(u);
                                                                 if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
                                                                  dis[m] = dis[u] - edge[v][m] + edge[u][v];
int Find(int u) {
                                                                  onstk[v] = 1;
 return u == fa[u] ? u : fa[u] = Find(fa[u]);
                                                                  stk.PB(v)
                                                                  if (SPFA(m)) return true;
                                                                  stk.pop_back();
int LCA(int x, int y, int n) {
 static int tk = 0; tk++;
                                                                  onstk[v] = 0;
 x = Find(x), y = Find(y);
 for (; ; swap(x, y)) {
  if (x != n) {
   if (v[x] == tk) return x;
                                                               onstk[u] = 0;
   v[x] = tk;
                                                               stk.pop_back();
                                                               return false;
   x = Find(pre[match[x]]);
                                                              int solve() {
                                                               // find a match
void Blossom(int x, int y, int 1) {
 while (Find(x) != 1) {
                                                               for (int i=0; i<n; i+=2){
  pre[x] = y, y = match[x];
                                                                match[i] = i+1;
   if (s[y] == 1) q.push(y), s[y] = 0;
                                                                match[i+1] = i;
  if (fa[x] == x) fa[x] = 1;
  if (fa[y] == y) fa[y] = 1;
                                                               while (true){
                                                                int found = 0;
  x = pre[y];
                                                                for (int i=0; i<n; i++)</pre>
                                                                 dis[i] = onstk[i] = 0;
bool Bfs(int r, int n) {
                                                                for (int i=0; i<n; i++){
 for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;
                                                                 stk.clear()
 while (!q.empty()) q.pop();
                                                                 if (!onstk[i] && SPFA(i)){
 q.push(r);
                                                                  found = 1;
                                                                  while (SZ(stk)>=2){
 s[r] = 0;
 while (!q.empty()) {
                                                                   int u = stk.back(); stk.pop_back();
                                                                   int v = stk.back(); stk.pop_back();
  int x = q.front(); q.pop();
   for (int u : g[x]) {
                                                                   match[u] = v;
   if (s[u] == -1)
                                                                   match[v] = u;
    pre[u] = x, s[u] = 1;
                                                                  }
     if (match[u] == n) {
      for (int a = u, b = x, last; b != n; a = last, b =
                                                                if (!found) break;
      pre[a])
       last = match[b], match[b] = a, match[a] = b;
                                                               int ret = 0;
      return true;
                                                               for (int i=0; i<n; i++)
                                                                ret += edge[i][match[i]];
    q.push(match[u]);
    s[match[u]] = 0;
                                                               return ret>>1;
    } else if (!s[u] && Find(u) != Find(x)) {
    int 1 = LCA(u, x, n);
Blossom(x, u, 1);
                                                             } graph;
                                                             4.5 Minimum Cost Circulation
     Blossom(u, x, 1);
                                                             struct Edge { int to, cap, rev, cost; };
  }
                                                             vector<Edge> g[kN];
                                                             int dist[kN], pv[kN], ed[kN];
 return false;
                                                             bool mark[kN];
                                                             int NegativeCycle(int n) {
int Solve(int n) {
                                                             memset(mark, false, sizeof(mark));
 int res = 0;
                                                              memset(dist, 0, sizeof(dist));
 for (int x = 0; x < n; ++x) {
                                                              int upd = -1;
                                                              for (int i = 0; i <= n; ++i) {</pre>
  if (match[x] == n) res += Bfs(x, n);
                                                               for (int j = 0; j < n; ++j) {
 return res;
                                                                int idx = 0;
                                                                for (auto &e : g[j]) {
}}
                                                                 if(e.cap > 0 && dist[e.to] > dist[j] + e.cost){
      Minimum Weight Matching (Clique version)
                                                                  dist[e.to] = dist[j] + e.cost;
struct Graph {
                                                                  pv[e.to] = j, ed[e.to] = idx;
 // 0-base (Perfect Match)
                                                                  if (i == n) {
                                                                   upd = j;
 int n, edge[MXN][MXN]
 int match[MXN], dis[MXN], onstk[MXN];
                                                                   while(!mark[upd])mark[upd]=1,upd=pv[upd];
 vector<int> stk;
                                                                   return upd;
 void init(int _n) {
  n = _n;
```

idx++;

```
}
return -1;
int Solve(int n) {
int rt = -1, ans = 0;
while ((rt = NegativeCycle(n)) >= 0) {
 memset(mark, false, sizeof(mark));
 vector<pair<int, int>> cyc;
while (!mark[rt]) {
   cyc.emplace_back(pv[rt], ed[rt]);
  mark[rt] = true;
   rt = pv[rt];
  reverse(cyc.begin(), cyc.end());
  int cap = kInf;
  for (auto &i : cyc)
  auto &e = g[i.first][i.second];
   cap = min(cap, e.cap);
  for (auto &i : cyc) {
   auto &e = g[i.first][i.second];
   e.cap -= cap;
   g[e.to][e.rev].cap += cap;
   ans += e.cost * cap;
  }
return ans:
```

4.6 Flow Models

- · Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}$.
 - 2. For each edge (x,y,l,u), connect x o y with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect v o T with capacity -in(v).
 - To maximize, connect $t\to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the
 - maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect t o s with capacity ∞ and let the flow from S to T be f'. If $f+f'
 eq \sum_{v\in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e \mbox{,}$ where f_e corresponds to the flow of edge e on the graph.
- ullet Construct minimum vertex cover from maximum matching M on bipartite graph(X, Y)
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise. 2. DFS from unmatched vertices in X.

 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect $x\to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y\to x$ with (cost,cap)=(-c,1)
 - 3. For each edge with c<0, sum these cost as K, then increase d(y)
 - by 1, decrease $d(\boldsymbol{x})$ by 1
 - 4. For each vertex v with d(v) > 0, connect $S \to v$ with (cost, cap) =(0, d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =(0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source s o v, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity
 - 5. For $v \in \mathit{G}$, connect it with sink $v \rightarrow t$ with capacity K + 2T - $(\sum_{e \in E(v)} w(e)) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u,v).
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to \bar{v} .
 - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$. 2. Create edge (u,v) with capacity w with w being the cost of choosing
 - u without choosing v.

- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x, t) with capacity c_x and create edge (s, y) with ca-
- 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

4.7 Dinic

```
template <typename Cap = int64_t>
class Dinic{
private:
  struct E{
     int to, rev;
     Cap cap;
  int n, st, ed;
  vector<vector<E>> G;
  vector<int> lv, idx;
  bool BFS() {
     lv.assign(n, -1);
     queue<int> bfs;
     bfs.push(st); lv[st] = 0;
     while (not bfs.empty()){
       int u = bfs.front(); bfs.pop();
       for (auto e: G[u]) {
         if (e.cap <= 0 or lv[e.to]!=-1) continue;</pre>
         bfs.push(e.to); lv[e.to] = lv[u] + 1;
       }
     return lv[ed] != -1;
  Cap DFS(int u, Cap f){
  if (u == ed) return f;
     Cap ret = 0;
     for(int &i = idx[u]; i < int(G[u].size()); ++i) {</pre>
       auto &e = G[u][i];
       if (e.cap <= 0 or lv[e.to]!=lv[u]+1) continue;</pre>
       Cap nf = DFS(e.to, min(f, e.cap));
ret += nf; e.cap -= nf; f -= nf;
       G[e.to][e.rev].cap += nf;
       if (f == 0) return ret;
     if (ret == 0) lv[u] = -1;
     return ret;
public:
  void init(int n_) { G.assign(n = n_, vector<E>()); }
  void add_edge(int u, int v, Cap c){
  G[u].push_back({v, int(G[v].size()), c});
     G[v].push_back({u, int(G[u].size())-1, 0});
  Cap max_flow(int st_, int ed_){
  st = st_, ed = ed_; Cap ret = 0;
     while (BFS()) {
       idx.assign(n, 0);
       Cap f = DFS(st, numeric_limits<Cap>::max());
       ret += f;
       if (f == 0) break;
     return ret;
```

Minimum Cost Maximum Flow 4.8

```
class MiniCostMaxiFlow{
 using Cap = int; using Wei = int64_t;
 using PCW = pair<Cap,Wei>;
 static constexpr Cap INF_CAP = 1 << 30;</pre>
 static constexpr Wei INF_WEI = 1LL<<60;</pre>
private:
 struct Edge{
  int to, back;
  Cap cap; Wei wei;
  Edge() {}
  Edge(int a,int b, Cap c, Wei d):
   to(a),back(b),cap(c),wei(d)
```

```
{}
                                                                    const int maxn = 500 + 5;
 };
                                                                    int w[maxn][maxn], g[maxn];
int ori, edd;
                                                                    bool v[maxn], del[maxn];
                                                                    void add_edge(int x, int y, int c) {
vector<vector<Edge>> G;
vector<int> fa, wh;
                                                                    w[x][y] += c; w[y][x] += c;
vector<bool> inq;
 vector<Wei> dis;
                                                                    pair<int, int> phase(int n) {
PCW SPFA(){
                                                                     memset(v, false, sizeof(v));
                                                                     memset(g, 0, sizeof(g));
int s = -1, t = -1;
  fill(inq.begin(),inq.end(),false);
  fill(dis.begin(), dis.end(), INF_WEI);
                                                                     while (true) {
  queue<int> qq; qq.push(ori);
  dis[ori] = 0;
                                                                      int c = -1;
                                                                      for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
  while(not qq.empty()){
   int u=qq.front();qq.pop();
   inq[u] = false;
                                                                        if (c == -1 \mid | g[i] > g[c]) c = i;
   for(int i=0;i<SZ(G[u]);++i){</pre>
                                                                      if (c == -1) break;
    Edge e=G[u][i];
                                                                      v[s = t, t = c] = true;
    int v=e.to; Wei d=e.wei;
                                                                      for (int i = 0; i < n; ++i) {
    if(e.cap <= 0 | |dis[v] <= dis[u] + d)
                                                                       if (del[i] || v[i]) continue;
    dis[v] = dis[u] + d;
                                                                       g[i] += w[c][i];
    fa[v] = u, wh[v] = i;
                                                                      }
    if (inq[v]) continue;
                                                                     return make_pair(s, t);
    qq.push(v);
    inq[v] = true;
                                                                    int mincut(int n) {
                                                                     int cut = 1e9;
  if(dis[edd]==INF_WEI) return {-1, -1};
                                                                     memset(del, false, sizeof(del));
  Cap mw=INF_CAP;
                                                                     for (int i = 0; i < n - 1; ++i) {
  for(int i=edd;i!=ori;i=fa[i])
                                                                      int s, t; tie(s, t) = phase(n);
   mw=min(mw,G[fa[i]][wh[i]].cap);
                                                                      del[t] = true; cut = min(cut, g[t]);
                                                                      for (int j = 0; j < n; ++j) {
  for (int i=edd;i!=ori;i=fa[i]){
                                                                       w[s][j] += w[t][j]; w[j][s] += w[j][t];
   auto &eg=G[fa[i]][wh[i]];
   eg.cap -= mw;
   G[eg.to][eg.back].cap+=mw;
                                                                     return cut;
  return {mw, dis[edd]};
public:
                                                                    5
                                                                          Math
void init(int a,int b,int n){
                                                                         Prime Table
  ori=a,edd=b;
                                                                    1002939109, 1020288887, 1028798297, 1038684299,\\
  G.clear();G.resize(n);
                                                                    1041211027, 1051762951, 1058585963, 1063020809,
  fa.resize(n);wh.resize(n);
                                                                    1147930723, 1172520109, 1183835981, 1187659051,
  inq.resize(n); dis.resize(n);
                                                                    1241251303, 1247184097, 1255940849, 1272759031,
                                                                    1287027493, 1288511629, 1294632499, 1312650799,
 void add_edge(int st, int ed, Cap c, Wei w){
                                                                    1868732623, 1884198443, 1884616807, 1885059541\\
  G[st].emplace_back(ed,SZ(G[ed]),c,w);
                                                                    \begin{array}{c} 1909942399, 1914471137, 1923951707, 1925453197, \\ 1979612177, 1980446837, 1989761941, 2007826547, \\ 2008033571, 2011186739, 2039465081, 2039728567, \end{array}
  G[ed].emplace_back(st,SZ(G[st])-1,0,-w);
                                                                    2093735719, 2116097521, 2123852629, 2140170259,
PCW solve(){
                                                                    3148478261, 3153064147, 3176351071, 3187523093,\\
  Cap cc=0; Wei ww=0;
                                                                    3196772239, 3201312913, 3203063977, 3204840059
  while(true){
                                                                    3210224309, 3213032591, 3217689851, 3218469083, 3219857533, 3231880427, 3235951699, 3273767923,
  PCW ret=SPFA();
                                                                    3276188869, 3277183181, 3282463507, 3285553889,
   if(ret.first==-1) break;
                                                                    3319309027, 3327005333, 3327574903, 3341387953,
   cc+=ret.first:
                                                                    3373293941, 3380077549, 3380892997, 3381118801\\
   ww+=ret.first * ret.second:
                                                                          \lfloor \frac{n}{i} \rfloor Enumeration
                                                                    T_0 = 1, T_{i+1} = \lfloor \frac{n}{\lfloor \frac{n}{T_i + 1} \rfloor} \rfloor
  return {cc,ww};
}
                                                                    5.3 ax+by=gcd
} mcmf;
                                                                    // ax+ny = 1, ax+ny == ax == 1 \pmod{n}
                                                                    void exgcd(lld x,lld y,lld &g,lld &a,lld &b) {
4.9
      GomoryHu Tree
                                                                     if (y == 0) g=x,a=1,b=0;
int g[maxn];
                                                                     else exgcd(y,x%y,g,b,a),b=(x/y)*a;
vector<edge> GomoryHu(int n){
vector<edge> rt;
                                                                    5.4
                                                                           Pollard Rho
 for(int i=1;i<=n;++i)g[i]=1;</pre>
for(int i=2;i<=n;++i){</pre>
                                                                    // does not work when n is prime
  int t=g[i]
                                                                    // return any non-trivial factor
                  // clear flows on all edge
  flow.reset();
                                                                    llu pollard_rho(llu n){
  rt.push_back({i,t,flow(i,t)});
                                                                     static auto f=[](llu x,llu k,llu m){
  flow.walk(i); // bfs points that connected to i (use
                                                                      return add(k,mul(x,x,m),m);
    edges not fully flow)
  for(int j=i+1;j<=n;++j){</pre>
                                                                     if (!(n&1)) return 2;
   if(g[j]==t && flow.connect(j))g[j]=i; // check if i
                                                                     mt19937 rnd(120821011);
    can reach j
                                                                     while(true){
  }
                                                                      llu y=2, yy=y, x=rnd()%n, t=1;
                                                                      for(llu sz=2;t==1;sz<<=1) {</pre>
return rt;
                                                                       for(llu i=0;i<sz;++i){</pre>
                                                                         if(t!=1)break;
```

yy=f(yy,x,n);

t=gcd(yy>y?yy-y:y-yy,n);

4.10 Global Min-Cut

```
}
y=yy;
}
if(t!=1&&t!=n) return t;
}
}
```

5.5 Pi Count (Linear Sieve)

```
static constexpr int N = 1000000 + 5;
11d pi[N];
vector<int> primes;
bool sieved[N];
11d cube_root(11d x){
 1ld s=cbrt(x-static_cast<long double>(0.1));
 while(s*s*s <= x) ++s;
 return s-1;
11d square_root(11d x){
lld s=sqrt(x-static_cast<long double>(0.1));
 while(s*s \ll x) ++s;
 return s-1;
void init(){
 primes.reserve(N)
 primes.push_back(1);
 for(int i=2;i<N;i++) {</pre>
  if(!sieved[i]) primes.push_back(i);
  pi[i] = !sieved[i] + pi[i-1];
  for(int p: primes) if(p > 1) {
   if(p * i >= N) break;
   sieved[p * i] = true;
}
   if(p % i == 0) break;
  }
11d phi(11d m, 11d n) {
static constexpr int MM = 80000, NN = 500;
static lld val[MM][NN];
 if(m<MM&&n<NN&&val[m][n])return val[m][n]-1;</pre>
 if(n == 0) return m;
 if(primes[n] >= m) return 1;
 lld ret = phi(m,n-1)-phi(m/primes[n],n-1);
 if(m<MM&&n<NN) val[m][n] = ret+1;</pre>
 return ret;
11d pi_count(11d);
11d P2(11d m, 11d n) {
 11d sm = square_root(m), ret = 0;
 for(lld i = n+1;primes[i]<=sm;i++)</pre>
  ret+=pi_count(m/primes[i])-pi_count(primes[i])+1;
 return ret;
11d pi_count(11d m) {
 if(m < N) return pi[m];</pre>
 11d n = pi_count(cube_root(m));
return phi(m, n) + n - 1 - P2(m, n);
```

5.6 Strling Number

5.6.1 First Kind

 $S_1(n,k)$ counts the number of permutations of n elements with k disjoint cycles.

$$S_1(n,k) = (n-1) \cdot S_1(n-1,k) + S_1(n-1,k-1)$$

$$x(x+1)\dots(x+n-1) = \sum_{k=0}^n S_1(n,k)x^k$$

$$g(x) = x(x+1)\dots(x+n-1) = \sum_{k=0}^n a_k x^k$$

$$\Rightarrow g(x+n) = \sum_{k=0}^n \frac{b_k}{(n-k)!} x^{n-k},$$

$$b_k = \sum_{i=0}^k ((n-i)! a_{n-i}) \cdot (\frac{n^{k-i}}{(k-i)!})$$

5.6.2 Second Kind

 $S_2(n,k)$ counts the number of ways to partition a set of n elements into k nonempty sets.

$$S_2(n,k) = S_2(n-1,k-1) + k \cdot S_2(n-1,k)$$

```
S_2(n,k) = \sum_{i=0}^k {k \choose i} i^n (-1)^{k-i} = \sum_{i=0}^k \frac{(-1)^i}{i!} \cdot \frac{(k-i)^n}{(k-i)!}
```

5.7 Range Sieve

```
const int MAX_SQRT_B = 50000;
const int MAX_L = 200000 + 5;

bool is_prime_small[MAX_SQRT_B];
bool is_prime[MAX_L];

void sieve(lld l, lld r){
   // [l, r)
   for(lld i=2;i*i<r;i++) is_prime_small[i] = true;
   for(lld i=1;i<r;i++) is_prime[i-1] = true;
   if(l=1) is_prime[0] = false;
   for(lld i=2;i*i<r;i++){
    if(!is_prime_small[i]) continue;
   for(lld j=i*i;j*j<r;j+=i) is_prime_small[j]=false;
   for(lld j=std::max(2LL, (l+i-1)/i)*i;j<r;j+=i)
        is_prime[j-l]=false;
   }
}</pre>
```

5.8 Miller Rabin

```
bool isprime(llu x){
 static llu magic[]={2,325,9375,28178,\
          450775,9780504,1795265022};
 static auto witn=[](llu a,llu u,llu n,int t)
 ->bool{
  if (!(a = mpow(a%n,u,n)))return 0;
  while(t--){
   1lu a2=mul(a,a,n);
   if(a2==1 && a!=1 && a!=n-1)
    return 1;
   a = a2;
  }
  return a!=1;
 if(x<2)return 0;</pre>
 if(!(x&1))return x==2;
 llu x1=x-1; int t=0;
 while(!(x1&1))x1>>=1,t++;
 for(llu m:magic)if(witn(m,x1,x,t))return 0;
 return 1:
```

5.9 Extended Euler

```
a^b \equiv \begin{cases} a^b \mod \varphi(m) + \varphi(m) & \text{if } (a,m) \neq 1 \land b \geq \varphi(m) \\ a^b \mod \varphi(m) & \text{otherwise} \end{cases} \pmod m
```

5.10 Gauss Elimination

```
void gauss(vector<vector<double>> &d) {
  int n = d.size(), m = d[0].size();
  for (int i = 0; i < m; ++i) {
    int p = -1;
    for (int j = i; j < n; ++j) {
      if (fabs(d[j][i]) < eps) continue;
      if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
    }
    if (p == -1) continue;
    for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
    for (int j = 0; j < n; ++j) {
      if (i == j) continue;
      double z = d[j][i] / d[i][i];
      for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
    }
}</pre>
```

5.11 Fast Fourier Transform

```
const int mod = 10000000007;
const int M1 = 985661441; // G = 3
const int M2 = 998244353;
const int M3 = 1004535809;
int superBigCRT(int64_t A, int64_t B, int64_t C) {
   static_assert (M1 <= M2 && M2 <= M3);
   constexpr int64_t r12 = modpow(M1, M2-2, M2);</pre>
```

```
constexpr int64_t r13 = modpow(M1, M3-2, M3);
                                                                 cplx b2 = (fb[i] - fb[j].conj()) * r4;
  constexpr int64_t r23 = modpow(M2, M3-2, M3);
                                                                 if (i != j) {
                                                                  cplx c1 = (fa[j] + fa[i].conj());
  constexpr int64_t M1M2 = 1LL * M1 * M2 % mod;
                                                                  cplx c2 = (fa[j] - fa[i].conj()) * r2;
cplx d1 = (fb[j] + fb[i].conj()) * r3;
  B = (B - A + M2) * r12 % M2;
 C = (C - A + M3) * r13 % M3;
 C = (C - B + M3) * r23 % M3
                                                                  cplx d2 = (fb[j] - fb[i].conj()) * r4;
  return (A + B * M1 + C * M1M2) % mod;
                                                                  fa[i] = c1 * d1 + c2 * d2 * r5;
                                                                  fb[i] = c1 * d2 + c2 * d1:
namespace fft {
                                                                 fa[j] = a1 * b1 + a2 * b2 * r5;
                                                                 fb[j] = a1 * b2 + a2 * b1;
using VI = vector<int>;
using VL = vector<long long>;
const double pi = acos(-1);
                                                                fft(fa, sz), fft(fb, sz);
cplx omega[maxn + 1];
                                                                vector<int> res(sz);
                                                                for (int i = 0; i < sz; ++i) {
void prefft() {
for (int i = 0; i <= maxn; i++)</pre>
                                                                 long long a = round(fa[i].re), b = round(fb[i].re),
  omega[i] = cplx(cos(2 * pi * j / maxn),
                                                                       c = round(fa[i].im);
                                                                 res[i] = (a+((b \% p) << 15)+((c \% p) << 30)) \% p;
     sin(2 * pi * j / maxn));
                                                                }
void fft(vector<cplx> &v, int n) {
                                                                return res;
 int z = __builtin_ctz(n) - 1;
                                                               }}
 for (int i = 0; i < n; ++i) {
                                                               5.12 Chinese Remainder
  int x = 0, j = 0;
  for (;(1 << j) < n;++j) x^=(i >> j & 1)<<(z - j);
                                                               lld crt(lld ans[], lld pri[], int n){
  if (x > i) swap(v[x], v[i]);
                                                                lld M = 1, ret = 0;
                                                                for(int i=0;i<n;i++) M *= pri[i];</pre>
 for (int s = 2; s <= n; s <<= 1) {
                                                                for(int i=0;i<n;i++)</pre>
  int z = s >> 1;
                                                                 lld iv = (gcd(M/pri[i],pri[i]).FF+pri[i])%pri[i];
  for (int i = 0; i < n; i += s) {
                                                                 ret += (ans[i]*(M/pri[i])%M * iv)%M;
   for (int k = 0; k < z; ++k) {
  cplx x = v[i + z + k] * omega[maxn / s * k];
                                                                 ret %= M;
    v[i + z + k] = v[i + k] - x;
                                                                return ret;
    v[i + k] = v[i + k] + x;
                                                               }
                                                               /*
                                                               Another:
                                                               x = a1 \% m1
                                                               x = a2 \% m2
void ifft(vector<cplx> &v, int n) {
                                                               g = gcd(m1, m2)
 fft(v, n);
                                                               assert((a1-a2)%g==0)
 reverse(v.begin() + 1, v.end());
                                                               [p, q] = exgcd(m2/g, m1/g)
 for (int i=0;i<n;++i) v[i] = v[i] * cplx(1. / n, 0);
                                                               return a2+m2*(p*(a1-a2)/g)
                                                               \theta \ll x \ll 1cm(m1, m2)
VL convolution(const VI &a, const VI &b) {
                                                               */
 // Should be able to handle N <= 10^5, C <= 10^4
 int sz = 1;
                                                                      Berlekamp Massey
 while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
                                                               // x: 1-base, p[]: 0-base
 vector<cplx> v(sz);
 for (int i = 0; i < sz; ++i) {
                                                               template<size t N>
                                                               vector<llf> BM(llf x[N], size_t n){
  double re = i < a.size() ? a[i] : 0;</pre>
                                                                size_t f[N]={0},t=0;11f d[N];
  double im = i < b.size() ? b[i] : 0;</pre>
                                                                vector<llf> p[N];
  v[i] = cplx(re, im);
                                                                for(size_t i=1,b=0;i<=n;++i) {</pre>
                                                                 for(size_t j=0;j<p[t].size();++j)
d[i]+=x[i-j-1]*p[t][j];</pre>
 fft(v, sz);
 for (int i = 0; i <= sz / 2; ++i) {
                                                                 if(abs(d[i]-=x[i])<=EPS)continue;</pre>
  int j = (sz - i) & (sz - 1);
                                                                 f[t]=i;if(!t){p[++t].resize(i);continue;}
 cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj())
                                                                 vector<llf> cur(i-f[b]-1);
    * cplx(0, -0.25);
                                                                 11f k=-d[i]/d[f[b]];cur.PB(-k);
  if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i
                                                                 for(size_t j=0;j<p[b].size();j++)
  cur.PB(p[b][j]*k);</pre>
    ].conj()) * cplx(\hat{0}, -0.25);
  v[i] = x;
                                                                 if(cur.size()<p[t].size())cur.resize(p[t].size());</pre>
                                                                 for(size_t j=0;j<p[t].size();j++)cur[j]+=p[t][j];</pre>
 ifft(v, sz);
                                                                 if(i-f[b]+p[b].size()>=p[t].size()) b=t;
 VL c(sz);
                                                                 p[++t]=cur;
 for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
 return c;
                                                                return p[t];
VI convolution_mod(const VI &a, const VI &b, int p) {
                                                               5.14 NTT
 while (sz + 1 < a.size() + b.size()) sz <<= 1;</pre>
 vector<cplx> fa(sz), fb(sz);
                                                               template <int mod, int G, int maxn>
 for (int i = 0; i < (int)a.size(); ++i)</pre>
                                                               struct NTT {
  fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
                                                                static_assert (maxn == (maxn & -maxn));
 for (int i = 0; i < (int)b.size(); ++i)</pre>
                                                                int roots[maxn];
  fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
                                                                NTT () {
 fft(fa, sz), fft(fb, sz);
                                                                 int r = modpow(G, (mod - 1) / maxn);
 double r = 0.25 / sz;
                                                                 for (int i = maxn >> 1; i; i >>= 1) {
 cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
for (int i = 0; i <= (sz >> 1); ++i) {
                                                                  roots[i] = 1;
                                                                  for (int j = 1; j < i; j++)
  int j = (sz - i) & (sz - 1);
                                                                   roots[i + j] = modmul(roots[i + j - 1], r);
 cplx a1 = (fa[i] + fa[j].conj());
cplx a2 = (fa[i] - fa[j].conj()) * r2;
                                                                  r = modmul(r, r);
  cplx b1 = (fb[i] + fb[j].conj()) * r3;
```

```
// n must be 2^k, and 0 \le F[i] < mod
                                                                 Poly Sqrt() const { // Jacobi(coef[0], P) = 1
 void inplace_ntt(int n, int F[], bool inv = false) {
                                                                  if (n()==1) return {QuadraticResidue(coef[0], P)};
  for (int i = 0, j = 0; i < n; i++) {
  if (i < j) swap(F[i], F[j]);</pre>
                                                                  Poly X = Poly(*this, (n()+1) / 2).Sqrt().isz(n());
                                                                  return X.iadd(Mul(X.Inv()).isz(n())).imul(P/2+1);
   for (int k = n > 1; (j^k < k; k > = 1);
                                                                 pair<Poly, Poly> DivMod(const Poly &rhs) const {
                                                                  // (rhs.)back() != 0
  for (int s = 1; s < n; s *= 2) {
   for (int i = 0; i < n; i += s * 2) {
                                                                  if (n() < rhs.n()) return {{0}, *this};</pre>
    for (int j = 0; j < s; j++) {
                                                                  const int _n = n() - rhs.n() + 1;
     int a = F[i+j];
                                                                  Poly X(rhs); X.irev().isz(_n);
     int b = modmul(F[i+j+s], roots[s+j]);
                                                                  Poly Y(*this); Y.irev().isz(_n);
     F[i+j] = modadd(a, b); // a + b
                                                                  Poly Q = Y.Mul(X.Inv()).isz(_n).irev();
                                                                  X = rhs.Mul(Q), Y = *this;
fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;</pre>
     F[i+j+s] = modsub(a, b); // a - b
                                                                  return {Q, Y.isz(max(1, rhs.n() - 1))};
   }
  if (inv) {
                                                                 Poly Dx() const {
                                                                 Poly ret(n() - 1);
   int invn = modinv(n);
                                                                  fi(0, ret.n()) ret[i] = (i + 1) * coef[i + 1] % P;
   for (int i = 0; i < n; i++)</pre>
    F[i] = modmul(F[i], invn);
                                                                  return ret.isz(max(1, ret.n()));
   reverse(F + 1, F + n);
                                                                 Poly Sx() const {
  Poly ret(n() + 1);
                                                                  fi(0, n()) ret[i + 1]=ntt.minv(i + 1)*coef[i] % P;
};
const int P=2013265921, root=31;
                                                                  return ret;
const int MAXN=1<<20;</pre>
NTT<P, root, MAXN> ntt;
                                                                 Poly _tmul(int nn, const Poly &rhs) const {
                                                                  Poly Y = Mul(rhs).isz(n() + nn - 1);
5.15 Polynomial Operations
                                                                  return Poly(Y.data() + n() - 1, nn);
using VL = vector<LL>
#define fi(s, n) for (int i=int(s); i<int(n); ++i)</pre>
                                                                 VL _eval(const VL &x, const auto up)const{
#define Fi(s, n) for (int i=int(n); i>int(s); --i)
                                                                  const int _n = (int)x.size();
int n2k(int n) {
                                                                  if (!_n) return {};
                                                                  vector<Poly> down(_n * 2);
int sz = 1; while (sz < n) sz <<= 1;</pre>
 return sz;
                                                                  down[1] = DivMod(up[1]).second;
                                                                  fi(2,_n*2) down[i]=down[i/2].DivMod(up[i]).second;
                                                                  /* down[1] = Poly(up[1]).irev().isz(n()).Inv().irev()
template<int MAXN, LL P, LL RT> // MAXN = 2^k
                                                                      _tmul(_n, *this)
struct Poly { // coefficients in [0, P)
 static NTT<MAXN, P, RT> ntt;
                                                                  fi(2, _n * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
                                                                    1, down[i / 2]); */
 VL coef;
 int n() const { return coef.size(); } // n()>=1
                                                                  VL y(_n);
                                                                  fi(0, _n) y[i] = down[_n + i][0];
 LL *data() { return coef.data(); }
 const LL *data() const { return coef.data(); }
                                                                  return y;
 LL &operator[](size_t i) { return coef[i]; }
                                                                 static vector<Poly> _tree1(const VL &x) {
 const LL &operator[](size_t i)const{return coef[i];}
                                                                  const int _n = (int)x.size();
 Poly(initializer_list<LL> a) : coef(a) { }
                                                                  vector<Poly> up(_n * 2);
 explicit Poly(int _n = 1) : coef(_n) { }
Poly(const LL *arr, int _n) : coef(arr, arr + _n) {}
Poly(const Poly &p, int _n) : coef(_n) {
                                                                  fi(0, _n) up[_n + i] = \{(x[i] ? P - x[i] : 0), 1\};
                                                                  Fi(0, _n-1) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
  copy_n(p.data(), min(p.n(), _n), data());
                                                                  return up;
 Poly& irev(){return reverse(data(),data()+n()),*this;}
                                                                 VL Eval(const VL&x)const{return _eval(x,_tree1(x));}
                                                                 static Poly Interpolate(const VL &x, const VL &y) {
 Poly& isz(int _n) { return coef.resize(_n), *this; }
 Poly& iadd(const Poly &rhs) { // n() == rhs.n()
                                                                  const int _n = (int)x.size();
                                                                  vector<Poly> up = _tree1(x), down(_n * 2);
VL z = up[1].Dx()._eval(x, up);
  fi(0, n()) if ((coef[i]+=rhs[i]) >= P)coef[i]-=P;
  return *this;
                                                                  fi(0, _n) z[i] = y[i] * ntt.minv(z[i]) % P;
                                                                  fi(0, _n) down[_n + i] = \{z[i]\};

Fi(0, _n-1) down[i] = down[i * 2].Mul(up[i * 2 + 1])
 Poly& imul(LL k) {
 fi(0, n()) coef[i] = coef[i] * k % P;
                                                                   .iadd(down[i * 2 + 1].Mul(up[i * 2]));
  return *this;
                                                                  return down[1];
 Poly Mul(const Poly &rhs) const {
                                                                 Poly Ln() const { // coef[0] == 1
  const int _n = n2k(n() + rhs.n() - 1);
 Poly X(*this, _n), Y(rhs, _n);
ntt(X.data(), _n), ntt(Y.data(),
fi(0, _n) X[i] = X[i] * Y[i] % P;
                                                                  return Dx().Mul(Inv()).Sx().isz(n());
                                                                 Poly Exp() const \{ // coef[0] == 0 \}
                                                                  if (n() == 1) return {1};
  ntt(X.data(), _n, true);
                                                                  Poly X = Poly(*this, (n() + 1)/2).Exp().isz(n());
Poly Y = X.Ln(); Y[0] = P - 1;
  return X.isz(n() + rhs.n() - 1);
                                                                  fi(0, n()) if((Y[i] = coef[i] - Y[i]) < 0)Y[i]+=P;
 Poly Inv() const { // coef[0] != 0
                                                                  return X.Mul(Y).isz(n());
 if (n() == 1) return {ntt.minv(coef[0])};
  const int _n = n2k(n() * 2);
  Poly Xi = Poly(*this, (n() + 1)/2).Inv().isz(_n);
                                                                 Poly Pow(const string &K) const {
  Poly Y(*this, _n);
                                                                  int nz = 0;
                                                                  while (nz < n() && !coef[nz]) ++nz;</pre>
  ntt(Xi.data(), _n), ntt(Y.data(), _n);
  fi(0, _n) {
Xi[i] *= (2 - Xi[i] * Y[i]) % P
                                                                  LL nk = 0, nk2 = 0;
                                                                  for (char c : K) {
                                                                  nk = (nk * 10 + c - '0') % P;
nk2 = nk2 * 10 + c - '0';
   if((Xi[i] \% = P) < 0) Xi[i] += P;
                                                                   if (nk2 * nz >= n()) return Poly(n());
  ntt(Xi.data(), _n, true);
  return Xi.isz(n());
                                                                   nk2 %= P - 1;
```

```
if (!nk && !nk2) return Poly({1}, n());
                                                                   if (y % g != 0) return -1;
  Poly X(data() + nz, n() - nz * nk2);
                                                                   t /= g, y /= g, M /= g;
  LL x0 = X[0]
                                                                   Int h = 0, gs = 1;
  return X.imul(ntt.minv(x0)).Ln().imul(nk).Exp()
                                                                   for (; h * h < M; ++h) gs = gs * x % M;
                                                                   unordered_map<Int, Int> bs;
   .imul(ntt.mpow(x0, nk2)).irev().isz(n()).irev();
                                                                   for (Int s = 0; s < h; bs[y] = ++s)
 Poly InvMod(int L) { // (to evaluate linear recursion)
                                                                     y = y * x % M;
 Poly R\{1, 0\}; // *this * R mod x^L = 1 (*this[0] ==
                                                                   for (Int s = 0; s < M; s += h) {
                                                                      t = t * gs % M;
  for (int level = 0; (1 << level) < L; ++level) {</pre>
                                                                      if (bs.count(t)) return c + s + h - bs[t];
   Poly 0 = R.Mul(Poly(data(), min(2 << level, n())));
   Poly Q(2 \ll level); Q[0] = 1;
                                                                   return -1;
   for (int j = (1 << level); j < (2 << level); ++j)
Q[j] = (P - O[j]) % P;</pre>
   R = R.Mul(Q).isz(4 << level);
                                                                 5.18
                                                                        FloorSum
  }
                                                                 // @param n `n < 2^32`
  return R.isz(L);
                                                                 // @param m `1 <= m < 2^32`
 \textbf{static} \  \, \textbf{LL} \  \, \textbf{LinearRecursion}(\textbf{const} \  \, \textbf{VL\&a}, \textbf{const} \  \, \textbf{VL\&c}, \textbf{LL} \  \, \textbf{n}) \, \{
                                                                  // @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
                                                                 llu floor_sum_unsigned(llu n, llu m, llu a, llu b) {
  // a_n = \sum_{j=0}^{n-j} a_{n-j}
  const int k = (int)a.size();
                                                                  llu ans = 0:
                                                                  while (true)
  assert((int)c.size() == k + 1);
                                                                   if (a >= m) {
  Poly C(k + 1), W(\{1\}, k), M = \{0, 1\};
                                                                     ans += n * (n - 1) / 2 * (a / m); a %= m;
  fi(1, k + 1) C[k - i] = c[i] ? P - c[i] : 0;
  C[k] = 1
                                                                   if (b >= m) {
  while (n) {
                                                                    ans += n * (b / m); b %= m;
   if (n % 2) W = W.Mul(M).DivMod(C).second;
   n /= 2, M = M.Mul(M).DivMod(C).second;
                                                                   llu y_max = a * n + b;
                                                                   if (y_max < m) break;</pre>
  LL ret = 0:
                                                                   // y_max < m * (n + 1)
  fi(0, k) ret = (ret + W[i] * a[i]) % P;
                                                                   // floor(y_max / m) <= n
  return ret;
                                                                   n = (11u)(y_max / m), b = (11u)(y_max % m);
                                                                   swap(m, a);
}:
#undef fi
                                                                  return ans;
#undef Fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
                                                                 11d floor_sum(11d n, 11d m, 11d a, 11d b) {
template<> decltype(Poly_t::ntt) Poly_t::ntt = {};
                                                                  11u ans = 0;
5.16 FWT
                                                                  if (a < 0) {
                                                                   11u \ a2 = (a \% m + m) \% m;
/* xor convolution:
* x = (x0,x1) , y = (y0,y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
                                                                   ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
* x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1)
* z = (1/2) * z''
                                                                  if (b < 0) {
                                                                   11u b2 = (b \% m + m) \% m;
                                                                   ans -= 1ULL * n * ((b2 - b) / m);
 * or convolution:
                                                                   b = b2:
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
                                                                  return ans + floor_sum_unsigned(n, m, a, b);
 * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */
const LL MOD = 1e9+7;
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
                                                                 5.19 Quadratic residue
 for( int d = 1 ; d < N ; d <<= 1 ) {
                                                                 struct S {
  int d2 = d << 1;
                                                                  int MOD, w;
  for( int s = 0 ; s < N ; s += d2 )
                                                                  int64_t x, y;
   for( int i = s , j = s+d ; i < s+d ; i++, j++ ){ LL ta = x[i] , tb = x[j];
                                                                  S(int m, int w_=-1, int64_t x_=1, int64_t y_=0)
                                                                    : MOD(m), w(w_), x(x_), y(y_) {}
    x[ i ] = ta+tb;
                                                                  S operator*(const S &rhs) const {
    x[ j ] = ta-tb;
    if( x[ i ] >= MOD ) x[ i ] -= MOD;
if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
                                                                   int w_{-} = w
                                                                   if (w<sub>_</sub> == -1) w<sub>_</sub> = rhs.w;
                                                                   assert(w_ != -1 and w_ == rhs.w);
                                                                   return { MOD, w_,
(x * rhs.x + y * rhs.y % MOD * w) % MOD,
 }
 if( inv )
                                                                     (x * rhs.y + y * rhs.x) % MOD };
  for( int i = 0 ; i < N ; i++ ) {</pre>
   x[`i ] *= inv(`N, MOD );
x[ i ] %= MOD;
                                                                 int get_root(int n, int P) {
  }
                                                                   if (P == 2 or n == 0) return n;
                                                                   if (qpow(n, (P - 1) / 2, P) != 1) return -1;
5.17 DiscreteLog
                                                                   auto check = [&](int x) {
                                                                   return qpow(x, (P - 1) / 2, P); };
if (check(n) == P-1) return -1;
template<typename Int>
Int BSGS(Int x, Int y, Int M) {
                                                                   int64_t a; int w; mt19937 rnd(7122);
  // x^? \equiv y (mod M)
  Int t = 1, c = 0, g = 1;
                                                                   do { a = rnd() % P;
                                                                     w = ((a * a - n) & P + P) & P;
  for (Int M_ = M; M_ > 0; M_ >>= 1)
                                                                   } while (check(w) != P - 1);
    g = g * x % M;
                                                                   return qpow(S(P, w, a, 1), (P + 1) / 2).x;
  for (g = gcd(g, M); t % g != 0; ++c) {
    if (t == y) return c;
    t = t * x % M;
                                                                 5.20 De-Bruijn
```

```
int res[maxn], aux[maxn], sz;
void db(int t, int p, int n, int k) {
if (t > n) {
  if (n \% p == 0)
   for (int i = 1; i <= p; ++i)
    res[sz++] = aux[i];
 } else {
  aux[t] = aux[t - p];
  db(t + 1, p, n, k);
for (int i = aux[t - p] + 1; i < k; ++i) {
   aux[t] = i;
   db(t + 1, t, n, k);
  }
 }
int de_bruijn(int k, int n) {
  // return cyclic string of len k^n s.t. every string
 // of len n using k char appears as a substring.
 if (k == 1) {
  res[0] = 0;
  return 1:
 for (int i = 0; i < k * n; i++) aux[i] = 0;
 sz = 0:
 db(1, 1, n, k);
 return sz;
```

5.21 Simplex Construction

Standard form: maximize $\sum_{1 \leq i \leq n} c_i x_i$ such that for all $1 \leq j \leq m$, $\sum_{1 < i < n} A_{ji} x_i \leq b_j$ and $x_i \geq 0$ for all $1 \leq i \leq n$.

- 1. In case of minimization, let $c_i' = -c_i$
- 2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- 3. $\sum_{1 \leq i \leq n} A_{ji} x_i = b_j$
 - $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$
 - $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

5.22 Simplex

```
namespace simplex {
// maximize c^Tx under Ax <= B</pre>
// return VD(n, -inf) if the solution doesn't exist
// return VD(n, +inf) if the solution is unbounded
using VD = vector<double>;
using VVD = vector<vector<double>>;
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
VVD d;
vector<int> p, q;
void pivot(int r, int s) {
double inv = 1.0 / d[r][s];
for (int i = 0; i < m + 2; ++i)
  for (int j = 0; j < n + 2; ++j)
  if (i != r && j != s)
   d[i][j] -= d[r][j] * d[i][s] * inv;
 for(int i=0;i<m+2;++i) if (i != r) d[i][s] *= -inv;
for(int j=0;j<n+2;++j) if (j != s) d[r][j] *= +inv;</pre>
d[r][s] = inv; swap(p[r], q[s]);
bool phase(int z) {
int x = m + z;
while (true) {
 int s = -1:
  for (int i = 0; i <= n; ++i) {
  if (!z && q[i] == -1) continue;
  if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
  if (d[x][s] > -eps) return true;
  int r = -1
  for (int i = 0; i < m; ++i) {
  if (d[i][s] < eps) continue;
if (r == -1 || \</pre>
   d[i][n+1]/d[i][s] < d[r][n+1]/d[r][s]) r = i;
  if (r == -1) return false;
 pivot(r, s);
```

```
VD solve(const VVD &a, const VD &b, const VD &c) {
m = b.size(), n = c.size();
d = VVD(m + 2, VD(n + 2));
 for (int i = 0; i < m; ++i)
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
 p.resize(m), q.resize(n + 1);
 for (int i = 0; i < m; ++i)</pre>
 p[i] = n + i, d[i][n] = -1, d[i][n + 1] = b[i];
for (int i = 0; i < n; ++i) q[i] = i,d[m][i] = -c[i];
 q[n] = -1, d[m + 1][n] = 1;
 int r = 0;
 for (int i = 1; i < m; ++i)</pre>
  if (d[i][n + 1] < d[r][n + 1]) r = i;
 if (d[r][n + 1] < -eps) {
  pivot(r, n)
  if (!phase(1) \mid \mid d[m + 1][n + 1] < -eps)
   return VD(n, -inf);
  for (int i = 0; i < m; ++i) if (p[i] == -1) {
   int`s = min_element(d[i].begin(), d[i].end() - 1)
        - d[i].begin();
   pivot(i, s);
 if (!phase(0)) return VD(n, inf);
 VD x(n);
 for (int i = 0; i < m; ++i)</pre>
  if (p[i] < n) x[p[i]] = d[i][n + 1];
 return x;
```

5.23 Charateristic Polynomial

```
vector<vector<int>>> Hessenberg(const vector<vector<int</pre>
    >> &A) {
 int N = A.size();
 vector<vector<int>> H = A;
 for (int i = 0; i < N - 2; ++i) {
  if (!H[i + 1][i]) {
   for (int j = i + 2; j < N; ++j) {
    if (H[j][i]) {
     for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j
     ][k])
     for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k]
    ][j]);
     break;
    }
  if (!H[i + 1][i]) continue;
  int val = fpow(H[i + 1][i], kP - 2);
  for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
   for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL
   * H[i + 1][k] * (kP - coef)) % kP;
for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i +
    1] + 1LL * H[k][j] * coef) % kP;
  }
 }
 return H;
vector<int> CharacteristicPoly(const vector<vector<int</pre>
    >> &A) {
 int N = A.size();
 auto H = Hessenberg(A);
 for (int i = 0; i < N; ++i) {
  for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
 vector<vector<int>> P(N + 1, vector<int>(N + 1));
 P[0][0] = 1;
 for (int i = 1; i <= N; ++i) {
  P[i][0] = 0;
  for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1][j]
    1];
  int val = 1;
  for (int j = i - 1; j >= 0; --j)
   int coef = 1LL * val * H[j][i
                                     _ i] % kP:
   for (int k = 0; k <= j; ++k) P[i][k] = (P[i][k] + 1
    LL * P[j][k] * coef) % kP;
   if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
```

return tie(a.x, a.y) < tie(b.x, b.y);});</pre>

```
if (N & 1) {
                                                              vector<PT> s(SZ(d)<<1);</pre>
  for (int i = 0; i \le N; ++i) P[N][i] = kP - P[N][i];
                                                              int o = 0;
                                                              for(auto p: d) {
                                                               while(o \ge 2 \& cross(p-s[o-2], s[o-1]-s[o-2]) <= 0)
return P[N];
}
                                                               0--:
                                                              s[o++] = p;
    Geometry
6
                                                              for(int i=SZ(d)-2, t = o+1;i>=0;i--){
    Basic Geometry
                                                               while(o>=t&&cross(d[i]-s[o-2],s[o-1]-s[o-2])<=0)</pre>
using coord_t = int;
                                                               s[o++] = d[i];
using Real = double;
using Point = std::complex<coord_t>;
                                                             s.resize(o-1);
int sgn(coord_t x) {
                                                              return s;
return (x > 0) - (x < 0);
coord_t dot(Point a, Point b) {
                                                            6.4 3D Convex Hull
return real(conj(a) * b);
                                                            // return the faces with pt indexes
                                                            int flag[MXN][MXN];
coord_t cross(Point a, Point b) {
return imag(conj(a) * b);
                                                            struct Point{
                                                             ld x,y,z;
int ori(Point a, Point b, Point c) {
                                                              Point operator * (const 1d &b) const {
                                                               return (Point) {x*b,y*b,z*b};}
return sgn(cross(b - a, c - a));
                                                             Point operator * (const Point &b) const {
                                                              return(Point) {y*b.z-b.y*z,z*b.x-b.z*x,x*b.y-b.x*y};
bool operator<(const Point &a, const Point &b) {</pre>
                                                              }
return real(a) != real(b)
  ? real(a) < real(b) : imag(a) < imag(b);</pre>
                                                            Point ver(Point a, Point b, Point c) {
                                                             return (b - a) * (c - a);
int argCmp(Point a, Point b) {
// -1 / 0 / 1 <-> < / == / > (atan2)
                                                             vector<Face> convex_hull_3D(const vector<Point> pt) {
int qa = (imag(a) == 0
                                                              int n = SZ(pt), ftop = 0
                                                              REP(i,n) REP(j,n) flag[i][j] = 0;
   ? (real(a) < 0 ? 3 : 1) : (imag(a) < 0 ? 0 : 2));
                                                              vector<Face> now;
 int qb = (imag(b) == 0
                                                              now.emplace_back(0,1,2);
   ? (real(b) < 0 ? 3 : 1) : (imag(b) < 0 ? 0 : 2));
if (qa != qb)
                                                              now.emplace_back(2,1,0)
                                                              for (int i=3; i<n; i++){
 return sgn(qa - qb);
 return sgn(cross(b, a));
                                                               ftop++; vector<Face> next;
                                                              REP(j, SZ(now)) {
  Face& f=now[j]; int ff = 0;
template <typename V> Real area(const V & pt) {
                                                                ld d=(pt[i]-pt[f.a]).dot(
coord_t ret = 0;
                                                                  ver(pt[f.a], pt[f.b], pt[f.c]));
for (int i = 1; i + 1 < (int)pt.size(); i++)</pre>
                                                                if (d <= 0) next.push_back(f);</pre>
 ret += cross(pt[i] - pt[0], pt[i+1] - pt[0]);
                                                                if (d > 0) ff=ftop;
 return ret / 2.0;
                                                                else if (d < 0) ff=-ftop;</pre>
                                                                flag[f.a][f.b]=flag[f.b][f.c]=flag[f.c][f.a]=ff;
6.2 Circle Class
                                                               REP(j, SZ(now)) {
struct Circle { Point o; Real r; };
                                                               Face& f=now[j]
                                                                if (flag[f.a][f.b] > 0 &&
vector<Real> intersectAngle(Circle a, Circle b) {
                                                                  flag[f.a][f.b] != flag[f.b][f.a])
Real d2 = norm(a.o - b.o)
                                                                 next.emplace_back(f.a,f.b,i);
if (norm(A.r - B.r) >= d2)
                                                                if (flag[f.b][f.c] > 0 &&
 if (A.r < B.r)
                                                                  flag[f.b][f.c] != flag[f.c][f.b])
   return {-PI, PI};
                                                                 next.emplace_back(f.b,f.c,i);
 else
                                                                if (flag[f.c][f.a] > 0 &&
   return {};
                                                                  flag[f.c][f.a] != flag[f.a][f.c])
 if (norm(A.r + B.r) <= d2) return {};</pre>
                                                                 next.emplace_back(f.c,f.a,i);
Real dis = hypot(A.x - B.x, A.y - B.y);
Real theta = atan2(B.y - A.y, B.x - A.x);
                                                              now=next;
Real phi = acos((A.r * A.r + d2 - B.r * B.r) /
   (2 * A.r * dis))
                                                              return now;
Real L = theta - phi, R = theta + phi;
while (L < -PI) L += PI * 2;
while (R > PI) R -= PI * 2;
                                                            6.5 2D Farthest Pair
return { L, R };
                                                            // stk is from convex hull
                                                            n = (int)(stk.size());
                                                            int pos = 1, ans = 0; stk.push_back(stk[0]);
vector<Point> intersectPoint(Circle a, Circle b) {
                                                             for(int i=0;i<n;i++) {</pre>
Real d=o.dis(aa.o);
                                                             while(abs(cross(stk[i+1]-stk[i],
if (d >= r+aa.r || d <= fabs(r-aa.r)) return {};
Real dt = (r*r - aa.r*aa.r)/d, d1 = (d+dt)/2;</pre>
                                                                stk[(pos+1)%n]-stk[i])) >
                                                                abs(cross(stk[i+1]-stk[i]
Point dir = (aa.o-o); dir /= d;
Point pcrs = dir*d1 + o;
                                                                stk[pos]-stk[i]))) pos = (pos+1)%n;
                                                              ans = max({ans, dis(stk[i], stk[pos]),
dt=sqrt(max(0.0L, r*r - d1*d1)), dir=dir.rot90();
return {pcrs + dir*dt, pcrs - dir*dt};
                                                               dis(stk[i+1], stk[pos])});
6.3 2D Convex Hull
                                                            6.6 2D Closest Pair
template<typename PT>
                                                            struct cmp_y {
vector<PT> buildConvexHull(vector<PT> d) {
                                                             bool operator()(const P& p, const P& q) const {
sort(ALL(d), [](const PT& a, const PT& b){
                                                               return p.y < q.y;</pre>
```

```
// NOTE: Point is complex<Real>
multiset<P, cmp_y> s;
                                                               // cross(pt-line.st, line.dir)<=0 <-> pt in half plane
void solve(P a[], int n) {
                                                               struct Line {
 sort(a, a + n, [](const P& p, const P& q) {
                                                                 Point st, ed;
  return tie(p.x, p.y) < tie(q.x, q.y);</pre>
                                                                 Point dir:
                                                                 Line (Point _s, Point _e)
 11f d = INF; int pt = 0;
                                                                  : st(_s), ed(_e), dir(_e - _s) {}
 for (int i = 0; i < n; ++i) {
                                                               };
  while (pt < i \text{ and } a[i].x - a[pt].x >= d)
   s.erase(s.find(a[pt++]));
                                                               bool operator<(const Line &lhs, const Line &rhs) {</pre>
                                                                 if (int cmp = argCmp(lhs.dir, rhs.dir))
  auto it = s.lower_bound(P(a[i].x, a[i].y - d));
  while (it != s.end() and it->y - a[i].y < d)</pre>
                                                                   return cmp == -1;
   d = min(d, dis(*(it++), a[i]));
                                                                 return ori(lhs.st, lhs.ed, rhs.st) < 0;</pre>
  s.insert(a[i]);
                                                               Point intersect(const Line &A, const Line &B) {
                                                                 Real t = cross(B.st - A.st, B.dir) /
                                                                  cross(A.dir, B.dir);
                                                                 return A.st + t * A.dir;
6.7 kD Closest Pair (3D ver.)
                                                               }
11f solve(vector<P> v) {
 shuffle(v.begin(), v.end(), mt19937());
unordered_map<lld, unordered_map<lld,</pre>
                                                               Real HPI(vector<Line> &lines) {
                                                                 sort(lines.begin(), lines.end());
  unordered_map<lld, int>>> m;
                                                                 deque<Line> que;
 llf d = dis(v[0], v[1]);
                                                                 deque<Point> pt;
 auto Idx = [&d] (11f x) -> 11d {
  return round(x * 2 / d) + 0.1;
                                                                 que.push_back(lines[0]);
                                                                 for (int i = 1; i < (int)lines.size(); i++) {</pre>
 auto rebuild_m = [&m, &v, &Idx](int k) {
                                                                   if (argCmp(lines[i].dir, lines[i-1].dir) == 0)
  m.clear();
  for (int i = 0; i < k; ++i)
m[Idx(v[i].x)][Idx(v[i].y)]</pre>
                                                               #define POP(L, R) \
                                                                   while (pt.size() > 0 \
    [Idx(v[i].z)] = i;
                                                                     && ori(L.st, L.ed, pt.back()) < 0) \
 }; rebuild_m(2);
                                                                   pt.pop_back(), que.pop_back(); \
while (pt.size() > 0 \
 for (size_t i = 2; i < v.size(); ++i) {</pre>
  const lld kx = Idx(v[i].x), ky = Idx(v[i].y),
                                                                     && ori(R.st, R.ed, pt.front()) < 0) \
     kz = Idx(v[i].z); bool found = false;
                                                                     pt.pop_front(), que.pop_front();
  for (int dx = -2; dx <= 2; ++dx) {
                                                                   POP(lines[i], lines[i])
   const 11d nx = dx + kx;
                                                                   pt.push_back(intersect(que.back(), lines[i]));
   if (m.find(nx) == m.end()) continue;
                                                                   que.push_back(lines[i]);
   auto& mm = m[nx];
   for (int dy = -2; dy <= 2; ++dy) {
                                                                 POP(que.front(), que.back())
    const 11d ny = dy + ky;
                                                                 if (que.size() <= 1 ||</pre>
    if (mm.find(ny) == mm.end()) continue;
                                                                   argCmp(que.front().dir, que.back().dir) == 0)
    auto& mmm = mm[ny];
                                                                   return 0;
    for (int dz = -2; dz <= 2; ++dz) {
                                                                 pt.push_back(intersect(que.front(), que.back()));
     const 11d nz = dz + kz;
                                                                 return area(pt);
     if (mmm.find(nz) == mmm.end()) continue;
     const int p = mmm[nz];
     if (dis(v[p], v[i]) < d) {
  d = dis(v[p], v[i]);</pre>
                                                               6.10 Minkowski sum
                                                               vector<pll> Minkowski(vector<pll> A, vector<pll> B) {
      found = true;
                                                                hull(A), hull(B);
                                                                vector<pll> C(1, A[0] + B[0]), s1, s2;
                                                                for(int i = 0; i < SZ(A); ++i)
                                                                 s1.pb(A[(i + 1) % SZ(A)] - A[i]);
                                                                for(int i = 0; i < SZ(B);</pre>
                                                                                           <u>i</u>++
  if (found) rebuild_m(i + 1);
                                                                 s2.pb(B[(i + 1) % SZ(B)] - B[i]);
  else m[kx][ky][kz] = i;
                                                                for(int p1 = 0, p2 = 0; p1 < SZ(A) \mid \mid p2 < SZ(B);)
                                                                 if (p2 >= SZ(B)
 return d;
                                                                   | | (p1 < SZ(A) \&\& cross(s1[p1], s2[p2]) >= 0))
                                                                  C.pb(C.back() + s1[p1++]);
6.8 Simulated Annealing
                                                                  C.pb(C.back() + s2[p2++]);
                                                                return hull(C), C;
11f anneal() {
 mt19937 rnd_engine( seed );
 uniform_real_distribution< llf > rnd( 0, 1 );
                                                               6.11 intersection of line and circle
 const llf dT = 0.001:
                                                               vector<pdd> line_interCircle(const pdd &p1,
 // Argument p
llf S_{cur} = calc(p), S_{best} = S_{cur};
for ( llf T = 2000 ; T > EPS ; T -= dT ) {
                                                                   const pdd &p2,const pdd &c,const double r){
                                                                pdd ft=foot(p1,p2,c),vec=p2-p1;
  // Modify p to p_prime
                                                                double dis=abs(c-ft);
 const llf S_prime = calc( p_prime );
                                                                if(fabs(dis-r)<eps) return vector<pdd>{ft};
                                                                if(dis>r) return {};
  const llf delta_c = S_prime - S_cur
                                                                vec=vec*sqrt(r*r-dis*dis)/abs(vec);
  11f prob = min( ( 11f ) 1, exp( -delta_c / T ) );
  if ( rnd( rnd_engine ) <= prob )</pre>
                                                                return vector<pdd>{ft+vec,ft-vec};
   S_cur = S_prime, p = p_prime;
  if ( S_prime < S_best ) // find min</pre>
                                                               6.12 intersection of polygon and circle
   S_best = S_prime, p_best = p_prime;
                                                               // Divides into multiple triangle, and sum up
 return S_best;
                                                               // test by HDU2892
                                                               const double PI=acos(-1);
```

double _area(pdd pa, pdd pb, double r){

if(abs(pa)<abs(pb)) swap(pa, pb);</pre>

cc.r = hypot(cc.o.x-a.x, cc.o.y-a.y);

return cc;

```
if(abs(pb)<eps) return 0;</pre>
 double S, h, theta;
 double a=abs(pb), b=abs(pa), c=abs(pb-pa);
                                                                template<typename P>
 double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
double cosC = dot(pa,pb) / a / b, C = acos(cosC);
                                                                Circle MinCircleCover(const vector<P>& pts){
                                                                 random_shuffle(pts.begin(), pts.end());
 if(a > r){
                                                                 Circle c = \{ pts[0], 0 \};
  S = (C/2)*r*r;
                                                                 for(int i=0;i<(int)pts.size();i++){</pre>
 h = a*b*sin(C)/c;
                                                                  if (dist(pts[i], c.o) <= c.r) continue;</pre>
                                                                  c = { pts[i], 0 };
for (int j = 0; j < i; j++) {
  if(dist(pts[j], c.o) <= c.r) continue;</pre>
  if (h < r && B < PI/2)
   S = (acos(h/r)*r*r - h*sqrt(r*r-h*h));
                                                                   c.o = (pts[i] + pts[j]) / 2;
 else if(b > r){
                                                                   c.r = dist(pts[i], c.o);
for (int k = 0; k < j; k++) {</pre>
  theta = PI - B - asin(sin(B)/r*a);
  S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
                                                                    if (dist(pts[k], c.o) <= c.r) continue;</pre>
 else S = .5*sin(C)*a*b;
                                                                    c = getCircum(pts[i], pts[j], pts[k]);
 return S;
                                                                  }
double area_poly_circle(const vector<pdd> poly,
                                                                 }
  const pdd &0,const double r){
                                                                 return c;
 double S=0;
 for(int i=0;i<SZ(poly);++i)</pre>
                                                                      KDTree (Nearest Point)
                                                                6.16
  S+=_area(poly[i]-0,poly[(i+1)%SZ(poly)]-0,r)
    *ori(0,poly[i],poly[(i+1)%SZ(poly)]);
                                                               const int MXN = 100005;
 return fabs(S);
                                                                struct KDTree {
                                                                 struct Node {
                                                                  int x,y,x1,y1,x2,y2;
6.13 intersection of two circle
                                                                  int id,f;
Node *L, *R;
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
                                                                 } tree[MXN], *root;
 pdd o1 = a.0, o2 = b.0;
                                                                 int n:
 double r1 = a.R, r2 = b.R, d2 = abs2(o1 - o2),
                                                                 LL dis2(int x1, int y1, int x2, int y2) {
  LL dx = x1-x2, dy = y1-y2;
     d = sqrt(d2);
 if(d < max(r1, r2) - min(r1, r2) \mid | d > r1 + r2)
                                                                  return dx*dx+dy*dy;
  return 0:
 pdd u = (o1 + o2) * 0.5
                                                                 static bool cmpx(Node& a, Node& b){return a.x<b.x;}</pre>
  + (o1 - o2) * ((r2 * r2 - r1 * r1) / (2 * d2));
                                                                 static bool cmpy(Node& a, Node& b){return a.y<b.y;}</pre>
 double A = sqrt((r1 + r2 + d) * (r1 - r2 + d)
                                                                 void init(vector<pair<int,int>> ip) {
     *(r1 + r2 - d) *(-r1 + r2 + d));
                                                                  n = ip.size();
 pdd v = pdd(o1.Y - o2.Y, -o1.X + o2.X) * A
                                                                  for (int i=0; i<n; i++) {
  / (2 * d2);
                                                                   tree[i].id = i;
 p1 = u + v, p2 = u - v;
                                                                   tree[i].x = ip[i].first;
 return 1;
                                                                   tree[i].y = ip[i].second;
                                                                  root = build_tree(0, n-1, 0);
      tangent line of two circle
vector<Line> go(const Cir& c1,
                                                                 Node* build_tree(int L, int R, int d) {
  const Cir& c2, int sign1){
                                                                  if (L>R) return nullptr
 // sign1 = 1 for outer tang, -1 for inter tang
                                                                  int M = (L+R)/2; tree[M].f = d%2;
 vector<Line> ret;
                                                                  nth_element(tree+L, tree+M, tree+R+1, d%2?cmpy:cmpx);
 double d_sq = norm2( c1.0 - c2.0 );
                                                                  tree[M].x1 = tree[M].x2 = tree[M].x;
 if( d_sq < eps ) return ret;</pre>
                                                                  tree[M].y1 = tree[M].y2 = tree[M].y;
 double d = sqrt( d_sq );
                                                                  tree[M].L = build_tree(L, M-1, d+1);
 Pt v = (c2.0 - c1.0) / d;
                                                                  if (tree[M].L) {
 double c = (c1.R - sign1 * c2.R) / d;
                                                                   tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
 if( c * c > 1 ) return ret;
 double h = sqrt( max( 0.0 , 1.0 - c * c ) );
                                                                   tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
 for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
                                                                   tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
 Pt n = \{ v.X * c - sign2 * h * v.Y ,
   v.Y * c + sign2 * h * v.X };
                                                                  tree[M].R = build_tree(M+1, R, d+1);
  Pt p1 = c1.0 + n * c1.R;
                                                                  if (tree[M].R) {
 Pt p2 = c2.0 + n * (c2.R * sign1);
                                                                   tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
 if( fabs( p1.X - p2.X ) < eps and</pre>
                                                                   tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
    fabs( p1.Y - p2.Y ) < eps )
                                                                   tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
   p2 = p1 + perp(c2.0 - c1.0);
  ret.push_back( { p1 , p2 } );
                                                                  }
                                                                  return tree+M;
 return ret;
                                                                 int touch(Node* r, int x, int y, LL d2){
                                                                  LL dis = sqrt(d2)+1;
6.15 Minimum Covering Circle
                                                                  if (x<r->x1-dis || x>r->x2+dis ||
template<typename P>
                                                                    y<r->y1-dis || y>r->y2+dis)
Circle getCircum(const P &a, const P &b, const P &c){
                                                                   return 0;
 Real a1 = a.x-b.x, b1 = a.y-b.y;
                                                                  return 1;
 Real c1 = (a.x+b.x)/2 * a1 + (a.y+b.y)/2 * b1;
 Real a2 = a.x-c.x, b2 = a.y-c.y;
                                                                 void nearest(Node* r,int x,int y,int &mID,LL &md2) {
                                                                  if (!r || !touch(r, x, y, md2)) return;
LL d2 = dis2(r->x, r->y, x, y);
 Real c2 = (a.x+c.x)/2 * a2 + (a.y+c.y)/2 * b2;
 Circle cc;
 cc.o.x = (c1*b2-b1*c2)/(a1*b2-b1*a2);
                                                                  if (d2 < md2 \mid | (d2 == md2 \&\& mID < r->id)) {
 cc.o.y = (a1*c2-c1*a2)/(a1*b2-b1*a2);
                                                                   mID = r -> id;
```

md2 = d2;

```
// search order depends on split dim
                                                                    bool neq = last < 0 || \</pre>
  if ((r->f == 0 && x < r->x) ||
                                                                     memcmp(s + sa[i], s + last,
(p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
    (r->f == 1 \&\& y < r->y)) {
   nearest(r->L, x, y, mID, md2);
nearest(r->R, x, y, mID, md2);
                                                                    ns[q[last = sa[i]]] = nmxz += neq;
  } else {
                                                                   sais(ns, nsa, p+nn, q+n, t+n, c+z, nn, nmxz+1);
                                                                   pre(sa, c, n, z);
   nearest(r->R, x, y, mID, md2);
                                                                   for (int i = nn - 1; i >= 0; --i)
   nearest(r->L, x, y, mID, md2);
                                                                    sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
                                                                   induce(sa, c, s, t, n, z);
 int query(int x, int y) {
  int id = 1029384756;
                                                                  void build(const string &s) {
 LL d2 = 102938475612345678LL;
                                                                  for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
                                                                   _s[(int)s.size()] = 0; // s shouldn't contain 0
  nearest(root, x, y, id, d2);
                                                                  sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for(int i = 0; i < (int)s.size(); ++i) sa[i]=sa[i+1];
for(int i = 0; i < (int)s.size(); ++i) rev[sa[i]]=i;</pre>
  return id;
} tree;
                                                                   int ind = 0; hi[0] = 0;
                                                                   for (int i = 0; i < (int)s.size(); ++i) {</pre>
     Stringology
                                                                    if (!rev[i]) {
                                                                     ind = 0;
7.1 Hash
                                                                     continue;
class Hash {
 private:
                                                                    while (i + ind < (int)s.size() && \</pre>
  static constexpr int P = 127, Q = 1051762951;
                                                                     s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
  vector<int> h, p;
                                                                    hi[rev[i]] = ind ? ind-- : 0;
 public:
  void init(const string &s){
                                                                 }}
   h.assign(s.size()+1, 0); p.resize(s.size()+1);
   for (size_t i = 0; i < s.size(); ++i)</pre>
                                                                  7.3 Suffix Automaton
    h[i + 1] = add(mul(h[i], P), s[i]);
                                                                  struct Node{
   generate(p.begin(), p.end(),[x=1,y=1,this]()
                                                                   Node *green, *edge[26];
     mutable{y=x;x=mul(x,P);return y;});
                                                                   int max_len;
                                                                   Node(const int _max_len)
  int query(int 1, int r){ // 1-base (1, r]
                                                                    : green(NULL), max_len(_max_len){
   return sub(h[r], mul(h[1], p[r-1]));}
                                                                    memset(edge, 0, sizeof(edge));
                                                                  } *ROOT, *LAST;
7.2 Suffix Array
                                                                  void Extend(const int c) {
                                                                   Node *cursor = LAST;
namespace sfxarray {
bool t[maxn * 2];
                                                                   LAST = new Node((LAST->max_len) + 1);
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2];
                                                                   for(;cursor&&!cursor->edge[c]; cursor=cursor->green)
                                                                    cursor->edge[c] = LAST;
int x[maxn], p[maxn], q[maxn * 2];
                                                                   if (!cursor)
                                                                    LAST->green = ROOT;
// sa[i]: sa[i]-th suffix is the \
                                                                   else {
// i-th lexigraphically smallest suffix.
                                                                    Node *potential_green = cursor->edge[c];
// hi[i]: longest common prefix \
// of suffix sa[i] and suffix sa[i - 1].
                                                                    if((potential_green->max_len)==(cursor->max_len+1))
void pre(int *sa, int *c, int n, int z) {
                                                                     LAST->green = potential_green;
memset(sa, 0, sizeof(int) * n);
                                                                    else {
 memcpy(x, c, sizeof(int) * z);
                                                                  //assert(potential_green->max_len>(cursor->max_len+1));
                                                                     Node *wish = new Node((cursor->max_len) + 1);
                                                                     for(;cursor && cursor->edge[c]==potential_green;
void induce(int *sa,int *c,int *s,bool *t,int n,int z){
memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i)
if (sa[i] && !t[sa[i] - 1])
                                                                        cursor = cursor->green)
                                                                      cursor->edge[c] = wish;
                                                                     for (int i = 0; i < 26; i++)
   sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
                                                                      wish->edge[i] = potential_green->edge[i];
                                                                     wish->green = potential_green->green;
 memcpy(x, c, sizeof(int) * z);
 for (int i = n - 1; i \ge 0; --i)
                                                                     potential_green->green = wish;
  if (sa[i] && t[sa[i] - 1])
                                                                     LAST->green = wish;
   sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q,
bool *t, int *c, int n, int z) {
bool uniq = t[n - 1] = true;
                                                                  char S[10000001], A[10000001];
 int nn=0, nmxz=-1, *nsa = sa+n, *ns=s+n, last=-1;
                                                                  int main(){
                                                                   scanf("%d%s", &N, S)
 memset(c, 0, sizeof(int) * z);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
                                                                   ROOT = LAST = new Node(0);
                                                                   for (int i = 0; S[i]; i++)
Extend(S[i] - 'a');
 if (uniq) {
                                                                   while (N--){
  for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
                                                                    scanf("%s", A);
  return;
                                                                    Node *cursor = ROOT;
 for (int i = n - 2; i \ge 0; --i)
                                                                    bool ans = true;
 t[i] = (s[i] = s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
                                                                    for (int i = 0; A[i]; i++){
 pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i)
if (t[i] && !t[i - 1])
                                                                     cursor = cursor->edge[A[i] - 'a'];
                                                                     if (!cursor) {
                                                                      ans = false;
   sa[--x[s[i]]] = p[q[i] = nn++] = i;
                                                                      break;
 induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)
  if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
                                                                    puts(ans ? "Yes" : "No");
```

```
return 0;
}
7.4 KMP
vector<int> kmp(const string &s) {
 vector<int> f(s.size(), θ);
 /* f[i] = length of the longest prefix
   (excluding s[0:i]) such that it coincides with the suffix of s[0:i] of the same length */
 /* i + 1 - f[i] is the length of the
   smallest recurring period of s[0:i] */
 int k = 0:
 for (int i = 1; i < (int)s.size(); ++i) {</pre>
 while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
  if (s[i] == s[k]) ++k;
  f[i] = k;
 }
 return f;
vector<int> search(const string &s, const string &t) {
 // return 0-indexed occurrence of t in s
 vector<int> f = kmp(t), r;
 for (int i = 0, k = 0; i < (int)s.size(); ++i) {
 while(k > 0 && (k==(int)t.size() \mid \mid s[i]!=t[k]))
   k = f[k - 1];
  if (s[i] == t[k]) ++k;
  if (k == (int)t.size()) r.push_back(i-t.size()+1);
 return res;
7.5 Z value
char s[MAXN];
int len,z[MAXN];
void Z_value() {
 int i,j,left,right;
 z[left=right=0]=len;
 for(i=1;i<len;i++)</pre>
  j=max(min(z[i-left], right-i),0);
  for(;i+j<len&&s[i+j]==s[j];j++);
  if(i+(z[i]=j)>right)right=i+z[left=i];
7.6 Manacher
int z[maxn];
int manacher(const string& s) {
  string t = ".";
 for(char c: s) t += c, t += '.';
 int 1 = 0, r = 0, ans = 0;
 for (int i = 1; i < t.length(); ++i) {</pre>
 z[i] = (r > i ? min(z[2 * 1 - i], r - i) : 1);
  while (i - z[i] \ge 0 \&\& i + z[i] < t.length()) {
   if(t[i - z[i]] == t[i + z[i]]) ++z[i];
   else break;
  if (i + z[i] > r) r = i + z[i], l = i;
 for(int i=1;i<t.length();++i) ans = max(ans, z[i]-1);</pre>
 return ans;
      Lexico Smallest Rotation
string mcp(string s){
 int n = s.length();
 s += s:
 int i=0, j=1;
 while (i<n && j<n){</pre>
 int k = 0;
  while (k < n \&\& s[i+k] == s[j+k]) k++;
```

if (s[i+k] <= s[j+k]) j += k+1;</pre>

else i += k+1:

if (i == j) j++;

int ans = i < n ? i : j;
return s.substr(ans, n);</pre>

```
7.8 BWT
```

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a'
 vector<int> v[ SIGMA ];
 void BWT(char* ori, char* res){
  // make ori -> ori + ori
  // then build suffix array
 void iBWT(char* ori, char* res){
  for( int i = 0 ; i < SIGMA ; i ++ )</pre>
   v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )</pre>
   v[ ori[i] - BASE ].push_back( i );
  vector<int> a:
  for( int i = 0 , ptr = 0
for( auto j : v[ i ] ){
                     ptr = 0 ; i < SIGMA ; i ++ )
    a.push_back( j );
    ori[ ptr ++ ] = BASE + i;
  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
  res[ i ] = ori[ a[ ptr ] ];</pre>
   ptr = a[ ptr ];
  res[ len ] = 0;
} bwt;
```

7.9 Palindromic Tree

```
struct palindromic_tree{
 struct node{
 int next[26],f,len;
  int cnt,num,st,ed;
 node(int l=0):f(0),len(1),cnt(0),num(0) {
  memset(next, 0, sizeof(next)); }
 vector<node> st;
 vector<char> s;
 int last, n;
 void init(){
 st.clear();s.clear();last=1; n=0;
  st.push_back(0);st.push_back(-1);
  st[0].f=1;s.push_back(-1); }
 int getFail(int x){
 while(s[n-st[x].len-1]!=s[n])x=st[x].f;
  return x;}
 void add(int c){
  s.push_back(c-='a'); ++n;
  int cur=getFail(last);
  if(!st[cur].next[c]){
   int now=st.size();
   st.push_back(st[cur].len+2);
   st[now].f=st[getFail(st[cur].f)].next[c];
   st[cur].next[c]=now;
   st[now].num=st[st[now].f].num+1;
  last=st[cur].next[c];
  ++st[last].cnt;}
 void dpcnt() {
 for (int i=st.size()-1; i >= 0; i--)
   st[st[i].f].cnt += st[i].cnt;
int size(){ return st.size()-2;}
} pt;
int main() {
 string s; cin >> s; pt.init();
 for (int i=0; i<SZ(s); i++) {</pre>
 int prvsz = pt.size(); pt.add(s[i]);
 if (prvsz != pt.size()) {
  int r = i, l = r - pt.st[pt.last].len + 1;
   // pal @ [1,r]: s.substr(1, r-1+1)
 }
 return 0;
```

8 Misc

8.1 Theorems

8.1.1 Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

8.1.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

8.1.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\ldots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

8.1.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$

holds for all $1 \le k \le n$.

8.1.5 Havel-Hakimi algorithm

find the vertex who has greatest degree unused, connect it with other greatest vertex.

8.1.6 Hall's marriage theorem

Let G be a finite bipartite graph with bipartite sets X and Y. For a subset W of X, let $N_G(W)$ denote the set of all vertices in Y adjacent to some element of W. Then there is an X-saturating matching iff $\forall W\subseteq X, |W|\le |N_G(W)|$

8.1.7 Euler's planar graph formula

$$V - E + F = C + 1, E \le 3V - 6$$
(?)

8.1.8 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$

8.1.9 Lucas's theorem

 $\binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p}, \text{ where } m=m_kp^k+m_{k-1}p^{k-1}+\cdots+m_1p+m_0, \\ \text{and } n=n_kp^k+n_{k-1}p^{k-1}+\cdots+n_1p+n_0.$

8.1.10 Matroid Intersection

Given matroids $M_1=(G,I_1),M_2=(G,I_2),$ find maximum $S\in I_1\cap I_2.$ For each iteration, build the directed graph and find a shortest path from s to t.

- $s \to x : S \sqcup \{x\} \in I_1$
- $x \to t : S \sqcup \{x\} \in I_2$
- $y \to x : S \setminus \{y\} \sqcup \{x\} \in I_1$ (y is in the unique circuit of $S \sqcup \{x\}$)
- $x \to y: S \setminus \{y\} \sqcup \{x\} \in I_2$ (y is in the unique circuit of $S \sqcup \{x\}$)

Alternate the path, and |S| will increase by 1. Let $R=\min(\mathrm{rank}(I_1),\mathrm{rank}(I_2)),N=|G|$. In each iteration, |E|=O(RN). For weighted case, assign weight -w(x) and w(x) to $x\in S$ and $x\notin S$, resp. Use Bellman-Ford to find the weighted shortest path. The maximum iteration of Bellman-Ford is 2R+1.

8.2 DP-opt Condition

8.2.1 totally monotone (concave/convex)

 $\begin{array}{l} \forall i < i', j < j', B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}$

8.2.2 monge condition (concave/convex)

 $\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}$

8.3 Convex 1D/1D DP

```
struct segment {
 int i, 1, r
 segment() {}
 segment(int a, int b, int c): i(a), l(b), r(c) {}
inline 1ld f(int 1, int r){return dp[1] + w(1+1, r);}
void solve() {
 dp[0] = 0;
 deque<segment> dq; dq.push_back(segment(0, 1, n));
 for (int i = 1; i <= n; ++i) {
  dp[i] = f(dq.front().i, i);</pre>
  while(dq.size()&&dq.front().r<i+1) dq.pop_front();</pre>
  dq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
  while (dq.size() &&
   f(i, dq.back().1) < f(dq.back().i, dq.back().1))
    dq.pop_back();
  if (dq.size())
   int d = 1 << 20, c = dq.back().1;</pre>
   while (d \gg 1) if (c + d \ll dq.back().r)
    if(f(i, c+d) > f(dq.back().i, c+d)) c += d;
   dq.back().r = c; seg.l = c + 1;
  if (seg.1 <= n) dq.push_back(seg);</pre>
 }
}
```

8.4 ConvexHull Optimization

```
struct Line {
 mutable int64_t a, b, p;
 bool operator<(const Line &rhs) const { return a < rhs</pre>
 bool operator<(int64_t x) const { return p < x; }</pre>
}:
struct DynamicHull : multiset<Line, less<>> {
 static const int64_t kInf = 1e18;
 bool Isect(iterator x, iterator y)
  auto Div = [](int64_t a, int64_t b) {
    return a / b - ((a ^ b) < 0 && a % b);
  if (y == end()) { x->p = kInf; return false; }
if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
  else x->p = Div(y->b - x->b, x->a - y->a);
  return x->p >= y->p;
 void Insert(int64_t a, int64_t b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
  while (Isect(y, z)) z = erase(z);
  if (x != begin() \&\& Isect(--x, y)) Isect(x, y = erase)
    (y));
  while ((y = x) != begin() && (--x)->p >= y->p) Isect(
    x, erase(y));
 int64_t Query(int64_t x) {
  auto 1 = *lower_bound(x);
  return 1.a * x + 1.b;
```

8.5 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
  return k;
}
```

8.6 Cactus Matching

```
vector<int> init_g[maxn],g[maxn*2];
int n,dfn[maxn],low[maxn],par[maxn],dfs_idx,bcc_id;
void tarjan(int u){
   dfn[u]=low[u]=++dfs_idx;
   for(int i=0;i<(int)init_g[u].size();i++){
    int v=init_g[u][i];</pre>
```

```
if(v==par[u]) continue;
  if(!dfn[v]){
                                                              par[1]=-1;
   par[v]=u;
                                                              tarjan(1);
                                                              dfs(1,-1);
   tarjan(v);
                                                              printf("%d\n", max(dp[1][0], dp[1][1]));
   low[u]=min(low[u],low[v]);
   if(dfn[u]<low[v]){</pre>
                                                              return 0;
   g[u].push_back(v);
    g[v].push_back(u);
                                                             8.7 DLX
  }else{
                                                             struct DLX {
   low[u]=min(low[u],dfn[v]);
                                                              const static int maxn=210;
   if(dfn[v]<dfn[u]){</pre>
                                                              const static int maxm=210;
    int temp_v=u;
                                                              const static int maxnode=210*210;
    bcc_id++;
                                                              int n, m, size, row[maxnode], col[maxnode];
    while(temp_v!=v){
                                                              int U[maxnode], D[maxnode], L[maxnode], R[maxnode];
     g[bcc_id+n].push_back(temp_v);
                                                              int H[maxn], S[maxm], ansd, ans[maxn];
                                                              void init(int _n, int _m) {
     g[temp_v].push_back(bcc_id+n);
                                                               n = _n, m = _m;
     temp_v=par[temp_v];
                                                               for(int i = 0; i <= m; ++i) {</pre>
    g[bcc_id+n].push_back(v);
                                                                S[i] = 0;
    g[v].push_back(bcc_id+n);
                                                                U[i] = D[i] = i;
    reverse(g[bcc_id+n].begin(),g[bcc_id+n].end());
                                                                L[i] = i-1, R[i] = i+1;
                                                               R[L[0] = size = m] = 0;
                                                               for(int i = 1; i <= n; ++i) H[i] = -1;
int dp[maxn][2],min_dp[2][2],tmp[2][2],tp[2];
                                                              void Link(int r, int c) {
void dfs(int u,int fa){
                                                               ++S[col[++size] = c];
if(u<=n){
                                                               row[size] = r; D[size] = D[c];
                                                               U[D[c]] = size; U[size] = c; D[c] = size;
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
                                                               if(H[r] < 0) H[r] = L[size] = R[size] = size;</pre>
   if(v==fa) continue;
                                                               else {
   dfs(v,u);
                                                                R[size] = R[H[r]];
   memset(tp,0x8f,sizeof tp);
                                                                L[R[H[r]]] = size;
   if(v<=n){
                                                                L[size] = H[r];
    tp[0]=dp[u][0]+max(dp[v][0],dp[v][1]);
                                                                R[H[r]] = size;
    tp[1]=max(
     dp[u][0]+dp[v][0]+1
     dp[u][1]+max(dp[v][0],dp[v][1])
                                                              void remove(int c) {
                                                               L[R[c]] = \dot{L}[c]; \hat{R}[\dot{L}[c]] = R[c];
   }else{
                                                               for(int i = D[c]; i != c; i = D[i])
                                                                for(int j = R[i]; j != i; j = R[j]) {
U[D[j]] = U[j];
    tp[0]=dp[u][0]+dp[v][0]
    tp[1]=max(dp[u][0]+dp[v][1],dp[u][1]+dp[v][0]);
                                                                 D[U[j]] = D[j];
   dp[u][0]=tp[0],dp[u][1]=tp[1];
                                                                  --S[col[i]]:
 }else{
  for(int i=0;i<(int)g[u].size();i++){</pre>
                                                              void resume(int c) {
   int v=g[u][i];
                                                               L[R[c]] = c; R[L[c]] = c;
   if(v==fa) continue;
                                                               for(int i = U[c]; i != c; i = U[i])
                                                                for(int j = L[i]; j != i; j = L[j]) {
   dfs(v,u);
                                                                 U[D[j]] = j;
 min_dp[0][0]=0;
                                                                 D[U[j]] =
  min_dp[1][1]=1;
                                                                 ++S[col[j]];
  min_dp[0][1]=min_dp[1][0]=-0x3f3f3f3f;
  for(int i=0;i<(int)g[u].size();i++){</pre>
   int v=g[u][i];
                                                              void dance(int d) {
   if(v==fa) continue;
                                                               if(d>=ansd) return;
   memset(tmp,0x8f,sizeof tmp);
                                                               if(R[0] == 0) {
   tmp[0][0]=max(
                                                                ansd = d;
   \min_{dp[0][0]+\max(dp[v][0],dp[v][1])}
                                                                return:
    \min_{dp[0][1]+dp[v][0]}
                                                               int c = R[0];
   tmp[0][1]=min_dp[0][0]+dp[v][0]+1;
                                                               for(int i = R[0]; i; i = R[i])
   tmp[1][0]=max(
                                                                if(S[i] < S[c]) c = i;
    \min_{dp[1][0]+\max(dp[v][0],dp[v][1])}
                                                               remove(c);
    min_dp[1][1]+dp[v][0]
                                                               for(int i = D[c]; i != c; i = D[i]) {
                                                                ans[d] = row[i]
   tmp[1][1]=min_dp[1][0]+dp[v][0]+1;
                                                                for(int j = R[i]; j != i; j = R[j])
   memcpy(min_dp,tmp,sizeof tmp);
                                                                 remove(col[j]);
                                                                dance(d+1);
                                                                for(int j = L[i]; j != i; j = L[j])
  dp[u][1]=max(min_dp[0][1],min_dp[1][0]);
  dp[u][0]=min_dp[0][0];
                                                                 resume(col[j]);
                                                               resume(c);
int main(){
int m,a,b;
scanf("%d%d",&n,&m)
                                                             } sol;
                                                                  Tree Knapsack
for(int i=0;i<m;i++){</pre>
 scanf("%d%d",&a,&b);
                                                             int dp[N][K];PII obj[N];
  init_g[a].push_back(b);
                                                             vector<int> G[N];
  init_g[b].push_back(a);
                                                             void dfs(int u, int mx){
```

```
for(int s: G[u]) {
  if(mx < obj[s].first) continue;</pre>
  for(int i=0;i<=mx-obj[s].FF;i++)</pre>
   dp[s][i] = dp[u][i];
  dfs(s, mx - obj[s].first);
  for(int i=obj[s].FF;i<=mx;i++)</pre>
   dp[u][i] = max(dp[u][i],
    dp[s][i - obj[s].FF] + obj[s].SS);
int main(){
 int n, k; cin >> n >> k;
 for(int i=1;i<=n;i++){</pre>
  int p; cin >> p;
  G[p].push_back(i);
 cin >> obj[i].FF >> obj[i].SS;
 dfs(0, k); int ans = 0;
 for(int i=0;i<=k;i++) ans = max(ans, dp[0][i]);</pre>
 cout << ans << '\n';
return 0;
8.9
      N Queens Problem
vector< int > solve( int n ) {
 // no solution when n=2, 3
 vector< int > ret;
 if ( n % 6 == 2 ) {
  for ( int i = 2 ; i <= n ; i += 2 )
  ret.push_back( i );
  ret.push_back( 3 ); ret.push_back( 1 );
for ( int i = 7 ; i <= n ; i += 2 )</pre>
   ret.push_back( i );
  ret.push_back( 5 );
 } else if ( n % 6 == 3 ) {
  for ( int i = 4 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
  ret.push_back( 2 );
  for ( int i = 5 ; i <= n ; i += 2 )
  ret.push_back( i );</pre>
  ret.push_back( 1 ); ret.push_back( 3 );
 } else {
 for ( int i = 2 ; i <= n ; i += 2 )
   ret.push_back( i );
  for ( int i = 1 ; i <= n ; i += 2 )
   ret.push_back( i );
return ret;
8.10 Aliens Optimization
long long Alien() {
 long long c = kInf;
 for (int d = 60; d >= 0; --d) {
  // cost can be negative, depending on the problem.
  if (c - (1LL << d) < 0) continue;</pre>
  long long ck = c - (1LL \ll d)
 pair<long long, int> r = check(ck);
if (r.second == k) return r.first - ck * k;
  if (r.second < k) c = ck;
 pair<long long, int> r = check(c);
 return r.first - c * k;
      To Check When Submit
8.11
  · Array out of bound.
  • long long / double cast.
```

- Initialization.
- Delete all debug code.
- Check sample testcase after simple modification.