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Basic

1.1 vimrc

```
se is nu ru et tgc sc hls cin cino+=j1 sw
      =4 sts=4 bs=2 mouse=a "encoding=utf
      -8 1s=2
Wextra -Wshadow -Wfatal-errors -
      Wconversion -fsanitize=address.
      undefined, float-divide-by-zero, float
-cast-overflow && echo success<CR>
map <leader>z <ESC>:w<CR>:!g++ "%" -o "%<
" -02 -g -std=gnu++20 && echo
      success<CR>
map <leader>i <ESC>:!./"%<"<CR>
map <leader>r <ESC>:!cat 01.in && echo " ---" && ./"%<" < 01.in<CR>
map <leader>l :%d<bar>0r ~/t.cpp<CR>
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d "[:space:]" \| md5sum \| cut -
let c_no_curly_error=1
```

1.2 Default code

```
#include <bits/stdc++.h>
 using namespace std;
using i64 = long long;
using ll = long long;
 #define SZ(v) (ll)((v).size())
 #define pb emplace_back
#define AI(i) begin(i), end(i)
#define X first
 #define Y second
 template<class T> bool chmin(T &a, T b) {
       return b < a && (a = b, true); }</pre>
 template<class T> bool chmax(T &a, T b) {
       return a < b && (a = b, true); }</pre>
#ifdef KEV
#define DE(args...) kout("[ " + string(#
    args) + " ] = ", args)
void kout() { cerr << endl; }</pre>
 template<class T, class ...U> void kout(T
       a, U ...b) { cerr << a << ' ', kout
      (b...); }
 template<class T> void debug(T l, T r) {
      while (l != r) cerr << *l << " \n"[</pre>
      next(l)==r], ++l; }
#else
#define DE(...) 0
 #define debug(...) 0
 #endif
 int main() {
  cin.tie(nullptr)->sync_with_stdio(false
   );
return 0;
|}
```

1.3 Fast Integer Input

```
char buf[1 << 16], *p1 = buf, *p2 = buf;</pre>
char get() {
  if (p1 == p2) {
    p1 = buf;
p2 = p1 + fread(buf, 1, sizeof(buf),
          stdin):
  if (p1 == p2)
  return -1;
return *p1++;
char readChar() {
  char c = get();
  while (isspace(c))
    c = get();
  return c;
int readInt() {
  int x = 0;
  char c = get();
  while (!isdigit(c))
    c = get();
  while (isdigit(c)) {
   x = 10 * x + c - '0';
```

```
c = get();
   return x;
i }
```

1.4 Fast Python Input

```
import sys, os, io
input = io.BytesIO(os.read(0, os.fstat(0)
    .st_size)).readline
```

1.5 Pragma optimization

```
#pragma GCC optimize("Ofast", "no-stack-
protector", "no-math-errno", "unroll
     protector",
      -loops")
#pragma GCC target("sse,sse2,sse3,ssse3,
      sse4, sse4.2, popcnt, abm, mmx, avx, tune=
      native, arch=core-avx2, tune=core-avx2
#pragma GCC ivdep
```

Flows, Matching

2.1 Flow

```
template <typename F>
struct Flow {
 static constexpr F INF = numeric_limits
      <F>::max() / 2;
  struct Edge {
    int to;
    Edge(int to, F cap): to(to), cap(cap
        ) {}
 int n;
vector<Edge> e;
 vector<vector<int>> adj;
  vector<int> cur, h;
  Flow(int n) : n(n), adj(n) {}
 bool bfs(int s, int t) {
  h.assign(n, -1);
    queue<int> q;
   h[s] = 0;
    q.push(s);
    while (!q.empty()) {
     int u = q.front();
      q.pop();
      for (int i : adj[u]) {
        auto [v, c] = e[i];
        if (c > 0 & h[v] == -1) {
          h[v] = h[u] + 1;
          if (v == t) { return true; }
          q.push(v);
     }
   }
    return false;
   dfs(int u, int t, F f) {
   if (u == t) { return f; }
    int j = adj[u][i];
      auto [v, c] = e[j];
if (c > 0 && h[v] == h[u] + 1) {
       F a = dfs(v, t, min(r, c));
        e[j].cap -= a;
       e[j ^ 1].cap += a;
        if (r == 0) { return f; }
     }
   return f - r;
 // can be bidirectional
 void addEdge(int u, int v, F cf = INF,
      F cb = 0) {
   adj[u].push_back(e.size()), e.
         emplace_back(v, cf);
    adj[v].push_back(e.size()), e.
        emplace_back(u, cb);
```

```
F maxFlow(int s, int t) {
     F ans = 0;
     while (bfs(s, t)) {
       cur.assign(n, 0);
       ans += dfs(s, t, INF);
    return ans;
  // do max flow first
  vector<int> minCut() {
    vector<int> res(n);
     for (int i = 0; i < n; i++) { res[i]</pre>
          = h[i] != -1; }
     return res;
| };
```

2.2

```
MCMF
template <class Flow, class Cost>
struct MinCostMaxFlow {
public:
  static constexpr Flow flowINF =
       numeric_limits<Flow>::max();
  static constexpr Cost costINF =
       numeric_limits<Cost>::max();
  MinCostMaxFlow() {}
  MinCostMaxFlow(int n) : n(n), g(n) {}
  int addEdge(int u, int v, Flow cap,
       Cost cost) {
    int m = int(pos.size());
    pos.push_back({u, int(g[u].size())});
    g[u].push_back({v, int(g[v].size()),
         cap, cost});
    g[v].push_back({u, int(g[u].size()) -
          1, 0, -cost});
  }
  struct edge {
    int u, v;
    Flow cap, flow;
    Cost cost;
  edge getEdge(int i) {
    auto _e = g[pos[i].first][pos[i].
        second];
    auto _re = g[_e.v][_e.rev];
    return {pos[i].first, _e.v, _e.cap +
         _re.cap, _re.cap, _e.cost};
  vector<edge> edges() {
    int m = int(pos.size());
    vector<edge> result(m);
    for (int i = 0; i < m; i++) { result[</pre>
         i] = getEdge(i); }
    return result;
  pair<Flow, Cost> maxFlow(int s, int t,
       Flow flow_limit = flowINF) {
       return slope(s, t, flow_limit).
       back(); }
  vector<pair<Flow, Cost>> slope(int s,
    int t, Flow flow_limit = flowINF)
    vector<Cost> dual(n, 0), dis(n);
    vector<int> pv(n), pe(n), vis(n);
auto dualRef = [&]() {
      fill(dis.begin(), dis.end(),
           costINF);
      fill(pv.begin(), pv.end(), -1);
      fill(pe.begin(), pe.end(), -1);
fill(vis.begin(), vis.end(), false)
      struct 0 {
        Cost key;
        int u:
        bool operator<(Q o) const {</pre>
              return key > o.key; }
      priority_queue<Q> h;
      dis[s] = 0;
h.push({0, s});
      while (!h.empty()) {
        int u = h.top().u;
        h.pop();
        if (vis[u]) { continue; }
        vis[u] = true;
```

```
if (u == t) { break; }
for (int i = 0; i < int(g[u].size</pre>
              ()); i++) {
           auto e = g[u][i];
           if (vis[e.v] | l e.cap == 0)
                continue;
           Cost cost = e.cost - dual[e.v]
                + dual[u];
           if (dis[e.v] - dis[u] > cost) {
             dis[e.v] = dis[u] + cost;
             pv[e.v] = u;
pe[e.v] = i;
             h.push({dis[e.v], e.v});
      if (!vis[t]) { return false; }
for (int v = 0; v < n; v++) {</pre>
         if (!vis[v]) continue;
         dual[v] -= dis[t] - dis[v];
      return true;
    Flow flow = 0;
Cost cost = 0, prevCost = -1;
    vector<pair<Flow, Cost>> result;
    result.push_back({flow, cost});
    while (flow < flow_limit) {</pre>
      if (!dualRef()) break;
      Flow c = flow_limit - flow;
       for (int v = t; v != s; v = pv[v])
         c = min(c, g[pv[v]][pe[v]].cap);
      for (int v = t; v != s; v = pv[v])
         auto& e = g[pv[v]][pe[v]];
         e.cap -= c;
        g[v][e.rev].cap += c;
      Cost d = -dual[s];
      flow += c;
cost += c * d;
       if (prevCost == d) { result.
            pop_back(); }
       result.push_back({flow, cost});
      prevCost = cost;
    return result;
private:
 int n:
  struct _edge {
    int v, rev;
    Flow cap;
    Cost cost;
  vector<pair<int, int>> pos;
  vector<vector<_edge>> g;
```

GomoryHu Tree

```
auto gomory(int n, vector<array<int, 3>>
  Flow<int, int> mf(n);
  for (auto [u, v, c] : e) { mf.addEdge(u
       , v, c, c); }
  vector<array<int, 3>> res;
  vector<int> p(n);
  for (int i = 1; i < n; i++) {</pre>
    for (int j = 0; j < int(e.size()); j
     ++) { mf.e[j << 1].cap = mf.e[j</pre>
           << 1 | 1].cap = e[j][2]; }
    int f = mf.maxFlow(i, p[i]);
    auto cut = mf.minCut();
     for (int j = i + 1; j < n; j++) { if
          (cut[i] == cut[j] && p[i] == p[j]) { p[j] = i; }}
    res.push_back({f, i, p[i]});
  return res;
```

```
Global Minimum Cut
```

|// 0(V ^ 3)

```
template <tvpename F>
struct GlobalMinCut {
  static constexpr int INF =
        numeric_limits<F>::max() / 2;
  vector<int> vis, wei;
  vector<vector<int>> adj;
  GlobalMinCut(int n) : n(n), vis(n), wei
        (n), adj(n, vector<int>(n)) {}
  void addEdge(int u, int v, int w){
    adj[u][v] += w;
    adj[v][u] += w;
  int solve() {
     int sz = n;
    int res = INF, x = -1, y = -1;
auto search = [&]() {
       fill(vis.begin(), vis.begin() + sz,
             0);
       fill(wei.begin(), wei.begin() + sz,
       0);
x = y = -1;
       int mx, cur;
       for (int i = 0; i < sz; i++) {
         mx = -1, cur' = 0;
         for (int j = 0; j < sz; j++) {
           if (wei[j] > mx) {
             mx = wei[j], cur = j;
         vis[cur] = 1, wei[cur] = -1;
         x = y;

y = cur;
         for (int j = 0; j < sz; j++) {
           if (!vis[j]) {
             wei[j] += adj[cur][j];
         }
       return mx;
     while (sz > 1) {
       res = min(res, search());
       for (int i = 0; i < sz; i++) {
         adj[x][i] += adj[y][i];
         adj[i][x] = adj[x][i];
       for (int i = 0; i < sz; i++) {
         adj[y][i] = adj[sz - 1][i];
adj[i][y] = adj[i][sz - 1];
      SZ--:
     return res;
};
```

Bipartite Matching

```
struct BipartiteMatching {
  int n, m;
  vector<vector<int>> adj;
  vector<int> 1, r, dis, cur;
  BipartiteMatching(int n, int m) : n(n),
        m(m), adj(n), l(n, -1), r(m, -1),
  dis(n), cur(n) {}
void addEdge(int u, int v) { adj[u].
       push_back(v); }
  void bfs() {
    vector<int> q;
    for (int u = 0; u < n; u++) {
  if (l[u] == -1) {
         q.push_back(u), dis[u] = 0;
        dis[u] = -1;
      }
    for (int i = 0; i < int(q.size()); i</pre>
         ++) {
       int u = q[i];
      for (auto v : adj[u]) {
         if (r[v] != -1 && dis[r[v]] ==
              -1) {
           dis[r[v]] = dis[u] + 1;
           q.push_back(r[v]);
        }
```

```
}
  bool dfs(int u) {
     for (int &i = cur[u]; i < int(adj[u].</pre>
          size()); i++) {
       int v = adj[u][i];
if (r[v] == -1 || dis[r[v]] == dis[
            u] + 1 && dfs(r[v])) {
         l[u] = v, r[v] = u;
         return true;
      }
    }
     return false;
  int maxMatching() {
     int match = 0;
     while (true) {
       bfs();
       fill(cur.begin(), cur.end(), 0);
       int cnt = 0;
       for (int u = 0; u < n; u++) {
         if (l[u] == -1) {
           cnt += dfs(u);
       if (cnt == 0) {
         break;
       match += cnt;
     return match;
  auto minVertexCover() {
     vector<int> L, R;
     for (int u = 0; u < n; u++) {
       if (dis[u] == -1) {
         L.push_back(u);
       } else if (l[u] != -1) {
         R.push_back(l[u]);
     return pair(L, R);
  }
};
```

GeneralMatching 2.6

```
struct GeneralMatchina {
  int n;
  vector<vector<int>> adj;
  vector<int> match;
  GeneralMatching(int n) : n(n), adj(n),
       match(n, -1) {}
  void addEdge(int u, int v) {
    adj[u].push_back(v);
    adj[v].push_back(u);
  int maxMatching() {
    vector<int> vis(n), link(n), f(n),
         dep(n);
    auto find = [&](int u) {
      while (f[u] != u) \{ u = f[u] = f[f[
           u]]; }
      return u;
    };
    auto lca = [&](int u, int v) {
      u = find(u);
      v = find(v);
      while (u != v) {
        if (dep[u] < dep[v]) { swap(u, v)
        u = find(link[match[u]]);
      return u;
    queue<int> q;
    auto blossom = [&](int u, int v, int
      while (find(u) != p) {
        link[u] = v;
        v = \overline{match[u]};
        if (vis[v] == 0) {
          vis[v] = 1;
          q.push(v);
        f[u] = f[v] = p;
        u = link[v];
```

```
}
      };
      auto augment = [&](int u) {
         while (!q.empty()) { q.pop(); }
          iota(f.begin(), f.end(), 0);
         fill(vis.begin(), vis.end(), -1);
q.push(u), vis[u] = 1, dep[u] = 0;
         while (!q.empty()){
             int u = q.front();
             q.pop();
            for (auto v : adj[u]) {
  if (vis[v] == -1) {
                  vis[v] = 0;
                  link[v] = u;
dep[v] = dep[u] + 1;
                  if (match[v] == -1) {
                     for (int x = v, y = u, tmp;

y!= -1; x = tmp, y =

x == -1 ? -1 : link[x
                           } ([
                        tmp = match[y], match[x]
                              = y, match[y] = x;
                     return true:
                  }
                  q.push(match[v]), vis[match[v
               ]] = 1, dep[match[v]] =
    dep[u] + 2;
} else if (vis[v] == 1 && find(
    v) != find(u)) {
                  int p = lca(u, v);
                  blossom(u, v, p), blossom(v,
                         u, p);
           }
         }
         return false;
      };
       int res = 0;
      for (int u = 0; u < n; ++u) { if (
    match[u] == -1) { res += augment</pre>
       return res;
   }
|};
```

Kuhn Munkres

```
| \ / \ | need perfect matching or not : w
 intialize with -INF / 0
template <typename Cost>
 struct KM {
   static constexpr Cost INF =
       numeric_limits<Cost>::max() / 2;
   vector<Cost> hl, hr, slk;
   vector<int> l, r, pre, vl, vr;
   queue<int> q;
   vector<vector<Cost>> w;
   KM(int n) : n(n), hl(n), hr(n), slk(n),
         l(n, -1), r(n, -1), pre(n), vl(n)
         , vr(n),
     w(n, vector<Cost>(n, -INF)) {}
   bool check(int x) {
     vl[x] = true;
if (l[x] != -1) {
       q.push(l[x]);
       return vr[l[x]] = true;
     while (x != -1) \{ swap(x, r[l[x] =
          pre[x]]); }
     return false;
   }
   void bfs(int s) {
     fill(slk.begin(), slk.end(), INF);
     fill(vl.begin(), vl.end(), false);
fill(vr.begin(), vr.end(), false);
     a = \{\};
     q.push(s);
     vr[s] = true;
     while (true) {
       Cost d;
       while (!q.empty()) {
          int y = q.front();
          q.pop();
          for (int x = 0; x < n; ++x) {
```

```
if (!vl[x] \&\& slk[x] >= (d = hl)
             [x] + hr[y] - w[x][y])) {
          pre[x] = y;
          if (d != 0) {
            slk[x] = d;
          } else if (!check(x)) {
            return;
          }
        }
     }
    d = INF:
    for (int x = 0; x < n; ++x) { if (!
         vl[x] && d > slk[x]) { d = slk}
         [x]; }}
    for (int x = 0; x < n; ++x) {
      if (vl[x]) {
        hl[x] += d;
      } else {
        slk[x] -= d;
      if (vr[x]) { hr[x] -= d; }
    for (int x = 0; x < n; ++x) { if (!
         vl[x] && !slk[x] && !check(x))
          { return; }}
 }
void addEdge(int u, int v, Cost x) { w[
    u][v] = max(w[u][v], x); }
Cost solve() {
  for (int i = 0; i < n; ++i) { hl[i] =
        *max_element(w[i].begin(), w[i
  ].end()); }
for (int i = 0; i < n; ++i) { bfs(i);
  Cost res = 0;
for (int i = 0; i < n; ++i) { res +=
      w[i][l[i]]; }
  return res:
```

Flow Models 2.8

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u - l.
 - 3. For each vertex v, denote by in(v)the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect t \rightarrow with capacity ∞ (skip this in circulation problem), and let fbe the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t
 - is the answer. To minimize, let f be the maximum G to T. Connect mum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f + f' \neq \sum_{v \in V, in(v) > 0} in(v),$ there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph
 - Redirect every edge: y → x if (x, y) ∈ M, x → y otherwise.
 DFS from unmatched vertices in X.

 - 3. $x \in X$ is chosen iff x is unvisited. 4. $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
 - 1. Consruct super source S and sink T2. For each edge (x, y, c), connect $x \rightarrow y$ with (cost, cap) = (c, 1) if c > y0, otherwise connect $y \rightarrow x$ with $(\cos t, cap) = (-c, 1)$

- 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
- 4. For each vertex v with d(v) > 0, connect $S \rightarrow v$ with (cost, cap) =(0, d(v))
- 5. For each vertex v with d(v) < 0, connect $v \to T$ with $(\cos t, cap) =$ (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C + K
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer T
 - 2. Construct a max flow model, let Kbe the sum of all weights
 - 3. Connect source $s \to v$, $v \in G$ with capacity K
 - 4. For each edge (u, v, w) in G, connect

 - flow f < K|V|
- · Minimum weight edge cover

 - where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$. Create edge (u,v) with capacity w
 - with w being the cost of choosing uwithout choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- \bullet 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + xyx'y')$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s, y) with capacity
- 2. Create edge (x, y) with capacity c_{xy} . 3. Create edge (x, y) and edge (x', y')with capacity $c_{xyx'y'}$.

Data Structure

<ext/pbds>

```
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef tree<int, null_type, std::less<
     int>, rb_tree_tag,
     tree_order_statistics_node_update>
     tree_set;
typedef cc_hash_table<int, int> umap;
typedef priority_queue<int> heap;
int main() {
  // rb tree
  tree set s:
  s.insert(71); s.insert(22);
  assert(*s.find_by_order(0) == 22);
       assert(*s.find_by_order(1) == 71);
  assert(s.order_of_key(22) == 0); assert
       (s.order_of_key(71) == 1);
  s.erase(22);
  assert(*s.find_by_order(0) == 71);
       assert(s.order_of_key(71) == 0);
  // mergable heap
  heap a, b; a.join(b);
  // persistant
  rope<char> r[2];
  r[1] = r[0];
```

};

```
std::string st = "abc"
                                                 int size(Treap *t) {
                                                                                                  array<Splay*, 2> ch = {nullptr, nullptr
   r[1].insert(0, st.c_str());
                                                     return t == nullptr ? 0 : t->sz;
                                                                                                  Splay* fa = nullptr;
   r[1].erase(1, 1);
                                                 void apply(Treap *t, i64 m, i64 b) {
   std::cout << r[1].substr(0, 2) << std::
                                                                                                  int sz = 1;
        endl;
                                                                                                  bool rev = false;
                                                     t->b += b;
   return 0:
                                                                                                  Splay() {}
                                                     t->val += m * size(t->lc) + b;
| }
                                                                                                  void applyRev(bool x) {
                                                                                                    if (x) {
                                                 void pull(Treap *t) {
 3.2 Li Chao Tree
                                                                                                       swap(ch[0], ch[1]);
                                                     t\rightarrow sz = size(t\rightarrow lc) + size(t\rightarrow rc) +
constexpr i64 INF = 4e18;
                                                                                                    }
 struct Line {
                                                                                                  }
                                                 void push(Treap *t) {
                                                                                                  void push() {
   i64 a, b;
                                                     if (t->lc != nullptr) {
   Line(): a(0), b(INF) {}
                                                                                                    for (auto k : ch) {
   Line(i64 a, i64 b) : a(a), b(b) {}
                                                         apply(t->lc, t->m, t->b);
                                                                                                      if (k) {
   i64 \ operator()(i64 \ x) \ \{ \ return \ a * x +
                                                                                                         k->applyRev(rev);
                                                     if (t->rc != nullptr) {
                                                         apply(t->rc, t->m, t->b + t->m *
};
// [, ) !!!!!!!!!!
struct Lichao {
                                                              (size(t->lc) + 1));
                                                                                                    rev = false;
                                                     t->m = t->b = 0:
  int n;
                                                                                                  void pull() {
   vector<int> vals;
                                                                                                     sz = 1;
                                                pair<Treap*, Treap*> split(Treap *t, int
   vector<Line> lines;
                                                                                                     for (auto k : ch) {
                                                     s) {
   Lichao() {}
                                                                                                       if (k) {
                                                     if (t == nullptr) { return {t, t}; }
   void init(const vector<int> &v) {
                                                                                                       }
                                                     push(t);
Treap *a, *b;
     n = v.size();
                                                                                                    }
     vals = v:
                                                     if (s <= size(t->lc)) {
    b = t;
     sort(vals.begin(), vals.end());
                                                                                                  int relation() { return this == fa->ch
     vals.erase(unique(vals.begin(), vals.
                                                                                                        \lceil 1 \rceil; \}
                                                         tie(a, b\rightarrow lc) = split(t\rightarrow lc, s);
          end()), vals.end());
                                                                                                  bool isRoot() { return !fa || fa->ch[0]
                                                     } else {
     lines.assign(4 * n, {});
                                                                                                         != this && fa->ch[1] != this; }
                                                         a = t;
                                                                                                  void rotate() {
                                                         tie(a->rc, b) = split(t->rc, s -
   int get(int x) { return lower_bound(
                                                                                                     Splay *p = fa;
                                                              size(t->lc) - 1);
        vals.begin(), vals.end(), x) -
vals.begin(); }
                                                                                                    bool x = !relation();
                                                                                                    p->ch[!x] = ch[x];
if (ch[x]) { ch[x]->fa = p; }
                                                     pull(t);
   void apply(Line p, int id, int l, int r
                                                     return {a, b};
        ) {
                                                                                                     fa = p -> fa;
     Line &q = lines[id];
                                                                                                     if (!p->isRoot()) { p->fa->ch[p->
                                                 Treap* merge(Treap *t1, Treap *t2) {
     if (p(vals[l]) < q(vals[l])) { swap(p
                                                     if (t1 == nullptr) { return t2; }
if (t2 == nullptr) { return t1; }
                                                                                                         relation()] = this; }
     , q); }
if (l + 1 == r) { return; }
                                                                                                    ch[x] = p;
p->fa = this;
                                                     push(t1), push(t2);
     int m = l + r >> 1;
                                                                                                    p->pull();
                                                     if (t1->w > t2->w) {
     if (p(vals[m]) < q(vals[m])) {</pre>
                                                         t1->rc = merge(t1->rc, t2);
                                                                                                  swap(p, q);
                                                         pull(t1);
       apply(p, id << 1, l, m);
                                                         return t1;
     } else {
                                                     } else {
       apply(p, id << 1 | 1, m, r);
                                                         t2->lc = merge(t1, t2->lc);
     }
                                                                                                         }
                                                         pull(t2)
                                                                                                     while (!s.empty()) {
                                                         return t2;
   void add(int ql, int qr, Line p) {
                                                                                                       s.back()->push();
     ql = get(ql), qr = get(qr);
                                                                                                       s.pop_back();
     auto go = [&](auto go, int id, int 1,
                                                 int rnk(Treap *t, i64 val) {
           int r) -> void {
                                                                                                    push();
                                                     int res = 0;
while (t != nullptr) {
       if (qr <= l || r <= ql) { return; }
                                                                                                     while (!isRoot()) {
       if (ql <= l && r <= qr) {
                                                                                                       if (!fa->isRoot()) {
                                                         push(t);
         apply(p, id, l, r);
                                                                                                         if (relation() == fa->relation())
                                                         if (val <= t->val) {
         return;
                                                             res += size(t->lc) + 1;
                                                                                                           fa->rotate();
                                                              t = t->rc;
       int m = l + r >> 1;
                                                                                                         } else {
                                                         } else {
       go(go, id << 1, l, m);
                                                             t = t->lc;
                                                                                                           rotate();
       go(go, id << 1 | 1, m, r);
                                                                                                         }
     go(go, 1, 0, n);
                                                     return res;
                                                                                                       rotate();
   i64 query(int p) {
                                                 Treap* join(Treap *t1, Treap *t2) {
                                                                                                    pull();
     p = get(p);
                                                     if (size(t1) > size(t2)) {
     auto go = [&](auto go, int id, int l,
                                                                                                  void access() {
  for (Splay *p = this, *q = nullptr; p
                                                         swap(t1, t2);
           int r) -> i64 {
       if (l + 1 == r) \{ return lines[id](
                                                     Treap *t = nullptr;
                                                                                                         ; q = p, p = p -> fa) {
                                                     while (t1 != nullptr) {
            vals[p]); }
                                                                                                       p->splay();
       int m = l + r >> 1;
                                                         auto [u1, v1] = split(t1, 1);
                                                                                                       p->ch[1] = q;
       return min(lines[id](vals[p]), p <</pre>
                                                         t1 = v1;
                                                                                                       p->pull();
            m ? go(go, id << 1, l, m) : go
(go, id << 1 | 1, m, r));
                                                         int r = rnk(t2, u1->val);
                                                         if (r > 0) {
                                                                                                     splay();
                                                             auto [u2, v2] = split(t2, r);
                                                              t = merge(t, u2);
     return go(go, 1, 0, n);
                                                                                                  void makeRoot() {
  }
                                                                                                    access();
                                                         }
|};
                                                                                                    applyRev(true);
                                                         t = merge(t, u1);
 3.3 Treap
                                                                                                  Splay* findRoot() {
                                                     t = merge(t, t2);
                                                                                                    access();
struct Treap {
                                                                                                     Splay *p = this;
     Treap *lc = nullptr, *rc = nullptr;
                                                                                                     while (p->ch[0]) \{ p = p->ch[0]; \}
     int sz = 1;
                                                                                                    p->splay();
     unsigned w = rng();
                                                 3.4 Link-Cut Tree
                                                                                                     return p;
     i64 m = 0, b = 0, val = 0;
```

| struct Splay {

```
friend void split(Splay *x, Splay *y) {
    x->makeRoot();
    y->access();
  // link if not connected
  friend void link(Splay *x, Splay *y) {
    x->makeRoot();
    if (y->findRoot() != x) {
      x->fa=y;
  }
  // delete edge if doesn't exist
  friend void cut(Splay *x, Splay *y) {
    split(x, y);
    if (x->fa == y \&\& !x->ch[1]) {
      x->fa = y->ch[0] = nullptr;
      x->pull();
    }
  bool connected(Splay *x, Splay *y) {
    return x->findRoot() == y->findRoot()
};
```

3.5 2D Segment Tree

```
int num[501][501], N, M; // input here
struct seg_2D {
  struct node {
    int data;
node *lc, *rc;
  } * root;
  node *merge(node *a, node *b, int 1,
    int r) {
node *p = new node;
    p->data = max(a->data, b->data);
    if (l == r) return p;
    int m = 1 + r >> 1;
    p->lc = merge(a->lc, b->lc, l, m);
    p->rc = merge(a->rc, b->rc, m+1, r)
    return p;
  node *build(int l, int r, int x) {
    node *p = new node;
    if (l == r) return p \rightarrow data = num[x][l
         ], p;
    int m = l + r >> 1;
    p->lc = build(l, m, x), p->rc = build
         (m + 1, r, x);
    p->data = max(p->lc->data, p->rc->
         data);
    return p;
  int query(int L, int R, int l, int r,
       node *p) {
    if (L <= l && R >= r) return p->data;
    int m = 1 + r >> 1, re = 0;
    if (L \ll m) re = query(L, R, l, m, p)
         ->lc);
    if (R > m)
      re = max(re, query(L, R, m + 1, r,
           p->rc));
    return re;
 }
};
struct seg_1D {
 struct node {
    seg_2D data;
node *lc, *rc;
  } * root;
 node *s_build(int l, int r) {
  node *p = new node;
    if (l == r)
      return p->data.root = p->data.build
          (1, M, l), p;
    int m = l + r >> 1;
    p->lc = s\_build(l, m), p->rc =
         s_build(m + 1, r);
    p->data.root = p->data.merge(
      p->lc->data.root, p->rc->data.root,
            1, M);
    return p;
  int s_query(int L, int R, int l, int r,
        node *p,
    int yl, int yr) {
```

```
if (L \le 1 \&\& R \ge r)
    return p->data.query(yl, yr, 1, M,
        p->data.root);
  int m = l + r >> 1, re = 0;
  if (L \ll m)
   re = s_query(L, R, l, m, p->lc, yl,
         yr);
 if (R > m)
   re = max(
     re, s_{query}(L, R, m + 1, r, p->rc
           , yl, yr));
void init() { root = s_build(1, N); }
int query(int xl, int xr, int yl, int
    yr) {
  return s_query(xl, xr, 1, N, root, yl
       , yr);
```

3.6 BIT Kth

3.7 Binary Index Tree

```
struct Binary_Index_Tree {
   int bit[MAXN + 1], lazy[MAXN + 1], n;
   int lb(int x) { return x & -x; }
   void init(int _n, int *data) {
     for (int i = 1, t; i <= n; ++i) {
       bit[i] = data[i], lazy[i] = 0, t =
            i - lb(i);
       for (int j = i - 1; j > t; j -= lb(
            j))
         bit[i] += bit[j];
   void suf_modify(int x, int v) {
     for (int t = x; t; t -= lb(t)) lazy[t
     for (int t = x + lb(x); t \&\& t <= n;
          t += lb(t))
       bit[t] += v * (x - t + lb(t));
   void modify(int x, int v) {
     for (; x; x \rightarrow b(x)) bit[x] += v;
   int query(int x) {
     int re = 0;
     for (int t = x; t; t -= lb(t))
re += lazy[t] * lb(t) + bit[t];
     for (int t = x + lb(x); t \&\& t <= n;
          t += lb(t)
       re += lazy[t] * (x - t + lb(t));
     return re;
};
```

3.8 DSU

```
| struct DSU {
    vector<int> arr;
    DSU(int n = 0): arr(n) {
        iota(ALL(arr), 0);
    }
    int boss(int x) {
        if (arr[x] == x) return x;
        return arr[x] = boss(arr[x]);
    }
    bool Union(int x, int y) {
        x = boss(x), y = boss(y);
        if (x == y) return 0;
        arr[y] = x;
        return 1;
    }
};
```

3.9 Interval Container

```
/* Add and remove intervals from a set of
      disjoint intervals.
 * Will merge the added interval with any
       overlapping intervals in the set
      when adding.
 * Intervals are [inclusive, exclusive).
set<pii>::iterator addInterval(set<pii>&
     is, int L, int R) {
  if (L == R) return is.end();
  auto it = is.lower_bound({L, R}),
       before = it;
  while (it != is.end() && it->X <= R) {</pre>
    R = max(R, it->Y);
    before = it = is.erase(it);
  if (it != is.begin() && (--it)->Y >= L)
    L = min(L, it->X);
    R = max(R, it->Y);
    is.erase(it);
  return is.insert(before, pii(L, R));
}
void removeInterval(set<pii>& is, int L,
     int R) {
  if (L == R) return;
  auto it = addInterval(is, L, R);
  auto r2 = it -> Y;
  if (it->X == L) is.erase(it);
  else (int&)it->Y = L;
  if (R != r2) is.emplace(R, r2);
```

3.10 KD Tree

```
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn],
     xr[maxn],
  yl[maxn], yr[maxn];
point p[maxn];
int build(int 1, int r, int dep = 0) {
   if (l == r) return -1;
  function<bool(const point &, const</pre>
       point &)> f =
     [dep](const point &a, const point &b)
       if (dep & 1) return a.x < b.x;</pre>
       else return a.y < b.y;</pre>
  int m = (l + r) >> 1;
nth_element(p + l, p + m, p + r, f);
  xl[m] = xr[m] = p[m].x;
   yl[m] = yr[m] = p[m].y;
   lc[m] = build(l, m, dep + 1);
  if (~lc[m]) {
    xl[m] = min(xl[m], xl[lc[m]]);
    xr[m] = max(xr[m], xr[lc[m]]);
    yl[m] = min(yl[m], yl[lc[m]]);
    yr[m] = max(yr[m], yr[lc[m]]);
  rc[m] = build(m + 1, r, dep + 1);
  if (~rc[m]) {
    xl[m] = min(xl[m], xl[rc[m]]);
     xr[m] = max(xr[m], xr[rc[m]]);
    yl[m] = min(yl[m], yl[rc[m]]);
    yr[m] = max(yr[m], yr[rc[m]]);
  }
  return m;
bool bound(const point &q, int o, long
     long d) {
  double ds = sqrt(d + 1.0);
  if (q.x < xl[o] - ds || q.x > xr[o] +
       ds II
     q.y < yl[o] - ds || q.y > yr[o] + ds)
     return false;
  return true;
long long dist(const point &a, const
     point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x)
     (a.y - b.y) * 111 * (a.y - b.y);
}
```

```
void dfs(
  const point &q, long long &d, int o,
   int dep = 0) {
   if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
   if (cd != 0) d = min(d, cd);
  if ((dep & 1) && q.x < p[o].x ||
     !(dep \& 1) \&\& q.y < p[o].y) {
     if (~lc[o]) dfs(q, d, lc[o], dep + 1)
    if (~rc[o]) dfs(q, d, rc[o], dep + 1)
    if (~rc[o]) dfs(q, d, rc[o], dep + 1)
    if (~lc[o]) dfs(q, d, lc[o], dep + 1)
  }
}
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i]</pre>
         = v[i];
  root = build(0, v.size());
long long nearest(const point &q) {
   long long res = 1e18;
   dfs(q, res, root);
   return res;
} // namespace kdt
```

3.11 Useful KD Tree

```
template <typename T, size_t kd> // kd
     ????????
class kd_tree {
public:
  struct point {
    T d[kd];
    inline T dist(const point &x) const {
      T ret = 0;
       for (size_t i = 0; i < kd; ++i)</pre>
         ret += std::abs(d[i] - x.d[i]);
      return ret;
    inline bool operator==(const point &p
       for (size_t i = 0; i < kd; ++i) {</pre>
         if (d[i] != p.d[i]) return 0;
      return 1;
    inline bool operator<(const point &b)</pre>
          const {
       return d[0] < b.d[0];</pre>
  };
private:
  struct node {
  node *1, *r;
    point pid;
    int s:
    node(const\ point\ \&p): l(0), r(0),
         pid(p), s(1) {}
    inline void up() {
      s = (1 ? 1 -> s : 0) + 1 + (r ? r -> s
            : 0);
 } * root;
const double alpha, loga;
const T INF; //????INF,?????
  int maxn;
  struct __cmp {
    int sort_id;
    inline bool operator()(
      const node *x, const node *y) const
      return operator()(x->pid, y->pid);
    inline bool operator()(
      const point &x, const point &y)
            const {
       if (x.d[sort_id] != y.d[sort_id])
         return x.d[sort_id] < y.d[sort_id</pre>
              ٦:
       for (size_t i = 0; i < kd; ++i) {</pre>
```

```
if (x.d[i] != y.d[i]) return x.d[
           i] < y.d[i];
    return 0;
  }
} cmp;
void clear(node *o) {
  if (!o) return;
  clear(o->1);
  clear(o->r);
  delete o;
inline int size(node *o) { return o ? o
     ->s : 0; }
std::vector<node *> A;
node *build(int k, int l, int r) {
  if (l > r) return 0;
if (k == kd) k = 0;
  int mid = (l + r) / 2;
  cmp.sort_id = k;
  std::nth_element(A.begin() + 1, A.
       begin() + mid,
  A.begin() + r + 1, cmp);
node *ret = A[mid];
  ret->l = build(k + 1, l, mid - 1);
  ret->r = build(k + 1, mid + 1, r);
  ret->up():
  return ret:
inline bool isbad(node *o) {
  return size(o->l) > alpha * o->s ||
    size(o->r) > alpha * o->s;
void flatten(node *u,
  typename std::vector<node *>::
       iterator &it) {
  if (!u) return;
  flatten(u->1, it);
  *it = u;
  flatten(u->r, ++it);
inline void rebuild(node *&u, int k) {
  if ((int)A.size() < u->s) A.resize(u
       ->s);
  typename std::vector<node *>::
       iterator it =
    A.begin();
  flatten(u, it);
  u = build(k, 0, u \rightarrow s - 1);
bool insert(
  node *&u, int k, const point &x, int
       dep) {
  if (!u) {
   u = new node(x);
    return dep <= 0;
  ++u->s;
  cmp.sort_id = k;
  if (insert(cmp(x, u->pid) ? u->l : u
        (k + 1) % kd, x, dep - 1)) {
    if (!isbad(u)) return 1;
   rebuild(u, k);
  return 0;
node *findmin(node *o, int k) {
  if (!o) return 0;
  if (cmp.sort_id == k)
    return o->l ? findmin(o->l, (k + 1)
         % kd) : o;
  node *l = findmin(o->l, (k + 1) % kd)
  node *r = findmin(o->r, (k + 1) % kd)
  if (l && !r) return cmp(l, o) ? l : o
  if (!l && r) return cmp(r, o) ? r : o
  if (!l && !r) return o:
  if (cmp(l, r)) return cmp(l, o) ? l :
  return cmp(r, o) ? r : o;
bool erase(node *&u, int k, const point
     &x) {
  if (!u) return 0;
```

```
if (u->pid == x) {
      if (u->r)
      else if (u->l) {
        u->r = u->1;
        u \rightarrow 1 = 0;
      } else {
        delete u;
        u = 0:
        return 1;
      --u->s;
      cmp.sort_id = k;
      u->pid = findmin(u->r, (k + 1) % kd
           )->pid;
      return erase(u->r, (k + 1) \% kd, u
           ->pid);
    cmp.sort_id = k;
    if (erase(cmp(x, u->pid) ? u->l : u->
          (k + 1) % kd, x)) {
      --u->s;
      return 1;
    } else return 0;
  inline T heuristic(const T h[]) const {
    T ret = 0;
    for (size_t i = 0; i < kd; ++i) ret</pre>
         += h[i];
    return ret;
  }
  int aM:
  std::priority_queue<std::pair<T, point</pre>
      >> pQ;
  void nearest(
    node *u, int k, const point &x, T *h,
          T &mndist) {
    if (u == 0 || heuristic(h) >= mndist)
          return;
    T dist = u->pid.dist(x), old = h[k];
    /*mndist=std::min(mndist,dist);*/
    if (dist < mndist) {</pre>
      pQ.push(std::make_pair(dist, u->pid
           ));
      if ((int)pQ.size() == qM + 1) {
        mndist = pQ.top().first, pQ.pop()
      }
    if (x.d[k] < u-pid.d[k]) {
      nearest(u->1, (k + 1) % kd, x, h,
           mndist):
      h[k] = std::abs(x.d[k] - u->pid.d[k]
           ]);
      nearest(u->r, (k + 1) % kd, x, h,
           mndist);
    } else {
      nearest(u->r, (k + 1) % kd, x, h,
           mndist);
      h[k] = std::abs(x.d[k] - u->pid.d[k]
           1):
      nearest(u->l, (k + 1) % kd, x, h,
           mndist);
    h[k] = old;
  std::vector<point> in_range;
  void range(
    node *u, int k, const point &mi,
         const point &ma) {
    if (!u) return;
    bool is = 1;
    for (int i = 0; i < kd; ++i)
      if (u->pid.d[i] < mi.d[i] ||</pre>
        ma.d[i] < u->pid.d[i]) {
        is = 0;
        break;
      }
    if (is) in_range.push_back(u->pid);
    if (mi.d[k] \leftarrow u \rightarrow pid.d[k])
      range(u->l, (k + 1) % kd, mi, ma);
    if (ma.d[k] >= u->pid.d[k])
      range(u->r, (k + 1) % kd, mi, ma);
public:
 kd_{tree}(const\ T\ \&INF,\ double\ a=0.75)
```

```
maxn(1) {}
   inline void clear() {
    clear(root), root = 0, maxn = 1;
  inline void build(int n, const point *p
       ) {
     clear(root), A.resize(maxn = n);
     for (int i = 0; i < n; ++i) A[i] =</pre>
         new node(p[i]);
    root = build(0, 0, n - 1);
  inline void insert(const point &x) {
    insert(root, 0, x, std::__lg(size(
         root)) / loga);
    if (root->s > maxn) maxn = root->s;
   inline bool erase(const point &p) {
    bool d = erase(root, 0, p);
    if (root && root->s < alpha * maxn)</pre>
         rebuild();
    return d;
  inline void rebuild() {
    if (root) rebuild(root, 0);
    maxn = root->s;
  inline T nearest(const point &x, int k)
    qM = k;
    T mndist = INF, h[kd] = \{\};
    nearest(root, 0, x, h, mndist);
mndist = pQ.top().first;
    pQ = std::priority_queue<std::pair<T,</pre>
          point>>();
    return mndist; /*???x?k?????*/
  inline const std::vector<point> &range(
    const point &mi, const point &ma) {
    in_range.clear();
    range(root, 0, mi, ma);
return in_range; /*???mi?ma????
         vector*/
   inline int size() { return root ? root
       ->s : 0; }
|};
```

3.12 Leftist Tree

```
struct node {
  ll v, data, sz, sum;
node *l, *r;
  node(ll k)
    : v(0), data(k), sz(1), l(0), r(0),
          sum(k) \{ \}
11 sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0;
node *merge(node *a, node *b) {
  if (!a | l !b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a \rightarrow r = merge(a \rightarrow r, b);
  if (V(a->r) > V(a->l)) swap(a->r, a->l)
  a -> v = V(a -> r) + 1, a -> sz = sz(a -> 1) +
       sz(a->r) + 1;
  a\rightarrow sum = sum(a\rightarrow 1) + sum(a\rightarrow r) + a\rightarrow
        data;
  return a;
void pop(node *&o) {
  node *tmp = o;
  o = merge(o \rightarrow l, o \rightarrow r);
  delete tmp;
```

3.13 Chmin Chmax

```
#include <algorithm>
#include <iostream>
using namespace std;
typedef long long ll;
```

```
const int MAXC = 200005;
const ll INF = 1e18;
struct node {
  ll sum;
  ll mx, mxcnt, smx;
  ll mi, micnt, smi;
  ll lazymax, lazymin, lazyadd;
  node(11 k = 0)
    : sum(k), mx(k), mxcnt(1), smx(-INF),
          mi(k),
      micnt(1), smi(INF), lazymax(-INF),
           lazymin(INF),
      lazyadd(0) {}
  node operator+(const node &a) const {
    node rt;
rt.sum = sum + a.sum;
    rt.mx = max(mx, a.mx);
    rt.mi = min(mi, a.mi);
    if (mx == a.mx) {
      rt.mxcnt = mxcnt + a.mxcnt;
      rt.smx = max(smx, a.smx);
    } else if (mx > a.mx) {
      rt.mxcnt = mxcnt;
      rt.smx = max(smx, a.mx);
    } else {
      rt.mxcnt = a.mxcnt;
      rt.smx = max(mx, a.smx);
    if (mi == a.mi) {
      rt.micnt = micnt + a.micnt;
      rt.smi = min(smi, a.smi);
    } else if (mi < a.mi) {</pre>
      rt.micnt = micnt;
      rt.smi = min(smi, a.mi);
    } else {
      rt.micnt = a.micnt;
      rt.smi = min(mi, a.smi);
    rt.lazymax = -INF;
    rt.lazymin = INF;
    rt.lazyadd = 0;
    return rt;
} seg[MAXC << 2];</pre>
11 a[MAXC];
void give_tag_min(int rt, ll t) {
  if (t >= seg[rt].mx) return;
  seg[rt].lazymin = t;
  seg[rt].lazymax = min(seg[rt].lazymax,
       t);
  seg[rt].sum -= seg[rt].mxcnt * (seg[rt
       ].mx - t);
  if (seg[rt].mx == seg[rt].smi) seg[rt].
       smi = t:
  if (seg[rt].mx == seg[rt].mi) seg[rt].
       mi = t;
  seg[rt].mx = t;
void give_tag_max(int rt, ll t) {
  if (t <= seg[rt].mi) return;</pre>
  seg[rt].lazymax = t;
  seg[rt].sum += seg[rt].micnt * (t - seg
       [rt].mi);
  if (seg[rt].mi == seg[rt].smx) seg[rt].
       smx = t;
  if (seg[rt].mi == seg[rt].mx) seg[rt].
       mx = t;
  seg[rt].mi = t;
}
void give_tag_add(int l, int r, int rt,
     ll t) {
  seg[rt].lazyadd += t;
  if (seg[rt].lazymax != -INF) seg[rt].
    lazymax += t;
  if (seg[rt].lazymin != INF) seg[rt].
       lazymin += t;
  seg[rt].mx += t;
  if (seg[rt].smx != -INF) seg[rt].smx +=
  seg[rt].mi += t;
  if (seg[rt].smi != INF) seg[rt].smi +=
  seg[rt].sum += (ll)(r - l + 1) * t;
```

```
void tag_down(int l, int r, int rt) {
  if (seg[rt].lazyadd != 0) {
    int mid = (l + r) >> 1;
    give_tag_add(l, mid, rt << 1, seg[rt
         ].lazyadd);
    give_tag_add(
      \label{eq:mid + 1, r, rt << 1 | 1, seg[rt].}
           lazyadd);
    seg[rt].lazyadd = 0;
  if (seg[rt].lazymin != INF) {
    give_tag_min(rt << 1, seg[rt].lazymin</pre>
    give_tag_min(rt \ll 1 | 1, seg[rt].
         lazymin);
    seg[rt].lazymin = INF;
  if (seg[rt].lazymax != -INF) {
    give_tag_max(rt << 1, seg[rt].lazymax</pre>
    give_tag_max(rt \ll 1 | 1, seg[rt].
         lazymax);
    seg[rt].lazymax = -INF;
}
void build(int l, int r, int rt) {
  if (l == r) return seg[rt] = node(a[l])
         void();
  int mid = (l + r) \gg 1;
  build(l, mid, rt << 1);
build(mid + 1, r, rt << 1 | 1);</pre>
  seg[rt] = seg[rt << 1] + seg[rt << 1]
       1];
void modifymax(
  int L, int R, int l, int r, int rt, ll
       t) {
  if (L <= l && R >= r && t < seg[rt].smi</pre>
    return give_tag_max(rt, t);
  if (l != r) tag_down(l, r, rt);
  int mid = (l + r) >> 1;
  if (L <= mid) modifymax(L, R, l, mid,</pre>
       rt << 1, t);
  if (R > mid)
    modifymax(L, R, mid + 1, r, rt << 1 |
          1, t);
  seg[rt] = seg[rt << 1] + seg[rt << 1]
       1];
void modifymin(
  int L, int R, int l, int r, int rt, ll
      t) {
  if (L \le 1 \&\& R \ge r \&\& t \ge seg[rt].smx
       )
    return give_tag_min(rt, t);
  if (l != r) tag_down(l, r, rt);
int mid = (l + r) >> 1;
  if (L <= mid) modifymin(L, R, 1, mid,</pre>
       rt << 1, t);
  if (R > mid)
    modifymin(L, R, mid + 1, r, rt << 1 |</pre>
          1, t);
  seg[rt] = seg[rt \ll 1] + seg[rt \ll 1]
       1];
void modifyadd(
  int L, int R, int l, int r, int rt, ll
       t) {
  if (L \ll 1 \&\& R \gg r)
    return give_tag_add(l, r, rt, t);
  if (l != r) tag_down(l, r, rt);
  int mid = (l + r) \gg 1;
  if (L <= mid) modifyadd(L, R, l, mid,</pre>
       rt << 1, t);
  if (R > mid)
    modifyadd(L, R, mid + 1, r, rt << 1 |
          1, t):
  seg[rt] = seg[rt << 1] + seg[rt << 1 |</pre>
       1];
```

```
}
ll query(int L, int R, int l, int r, int
     rt) {
   if (L \ll l \& R \gg r) return seg[rt].
       sum;
  if (l != r) tag_down(l, r, rt);
  int mid = (l + r) \gg 1;
  if (R <= mid) return query(L, R, l, mid</pre>
        , rt << 1);
  if (L > mid)
    return query(L, R, mid + 1, r, rt <<</pre>
         1 | 1);
  return query(L, R, l, mid, rt << 1) +</pre>
    query(L, R, mid + 1, r, rt << 1 | 1);
int main() {
  ios::sync_with_stdio(0), cin.tie(0);
   int n, m;
  cin >> n >> m;
  for (int i = 1; i <= n; ++i) cin >> a[i
  build(1, n, 1);
  while (m--) {
    int k, x, y;
    ll t;
    cin >> k >> x >> y, ++x;
    if (k == 0) cin >> t, modifymin(x, y,
          1, n, 1, t);
     else if (k == 1)
      cin >> t, modifymax(x, y, 1, n, 1,
           t);
    else if (k == 2)
       cin >> t, modifyadd(x, y, 1, n, 1,
           t);
    else cout << query(x, y, 1, n, 1) <<
  }
|}
```

3.14 Segment Tree

```
struct Segment_Tree {
  struct node {
    int data, lazy;
node *1, *r;
node() : data(0), lazy(0), l(0), r(0)
          {}
    void up() {
      if (1) data = max(1->data, r->data)
    void down() {
      if (1) {
        l->data += lazy, l->lazy += lazy;
r->data += lazy, r->lazy += lazy;
      lazy = 0;
 }
} * root;
  int l, r;
  node *build(int l, int r, int *data) {
    node *p = new node();
    if (l == r) return p->data = data[l],
          р;
    int m = (l + r) / 2;
    p->l = build(l, m, data),
    p->r = build(m + 1, r, data);
    return p->up(), p;
 }
  void s_modify(
    int L, int R, int l, int r, node *p,
         int x) {
    if (r < L || l > R) return;
    p->down();
    if (L <= 1 && R >= r)
      return p->data += x, p->lazy += x,
           void();
    int m = (l + r) / 2;
    s_{modify(L, R, l, m, p->l, x)};
    s_{modify}(L, R, m + 1, r, p->r, x);
    p->up();
  int s_query(int L, int R, int l, int r,
        node *p) {
    p->down();
    if (L <= 1 && R >= r) return p->data;
```

3.15 Smart Pointer

```
#ifndef REFERENCE_POINTER
#define REFERENCE_POINTER
template <typename T> struct _RefCounter
  T data;
  int ref;
  _RefCounter(const T &d = 0) : data(d),
       ref(0) {}
témplate <typename T> struct
     reference_pointer {
   RefCounter<T> *p;
  T *operator->() { return &p->data; }
T &operator*() { return p->data; }
  operator _RefCounter<T> *() { return p;
  reference_pointer &operator=(
    const reference_pointer &t) {
    if (p && !--p->ref) delete p;
    p = t.p;
    p && ++p->ref;
    return *this;
  reference_pointer(_RefCounter<T> *t =
       0) : p(t) {
    p && ++p->ref;
  reference_pointer(const
       reference_pointer &t)
    : p(t.p) {
    p && ++p->ref;
  }
  ~reference_pointer() {
    if (p && !--p->ref) delete p;
template <typename T>
inline reference_pointer<T> new_reference
  const T &nd) {
  return reference_pointer<T>(new
       _RefCounter<T>(nd));
}
#endif
// note:
reference_pointer<int> a;
a = new_reference(5);
a = new_reference<int>(5);
a = new_reference((int)5);
reference_pointer<int> b = a;
struct P {
  int a, b;
  P(int _a, int _b) : a(_a), b(_b) {}
} p(2, 3);
reference_pointer<P> a;
c = new_reference(P(1, 2));
c = new\_reference < P > (P(1, 2));
c = new_reference(p);
3.16 Sparse Table
```

```
| struct Sparse_table {
| int st[__lg(MAXN) + 1][MAXN], n;
| void init(int _n, int *data) {
```

3.17 Discrete Trick

```
vector<int> val;
// build
sort(ALL(val)), val.resize(unique(ALL(val)) - val.begin());
// index of x
upper_bound(ALL(val), x) - val.begin();
// max idx <= x
upper_bound(ALL(val), x) - val.begin();
// max idx < x
lower_bound(ALL(val), x) - val.begin();</pre>
```

3.18 Min Heap

```
template<class T, class Info>
struct min_heap {
  priority_queue<pair<T, Info>, vector<</pre>
       pair<T, Info>>, greater<pair<T,</pre>
       Info>>> pq;
  T lazy = 0;
  void push(pair<T, Info> v) {
    pq.emplace(v.X - lazy, v.Y);
  pair<T, Info> top() {
    return make_pair(pq.top().X + lazy,
         pq.top().Y);
   void join(min_heap &rgt) {
    if (SZ(pq) < SZ(rgt.pq)) {</pre>
       swap(pq, rgt.pq);
       swap(lazy, rgt.lazy);
    while (!rgt.pq.empty()) {
       push(rgt.top());
       rgt.pop();
    }
  void pop() {
    pq.pop();
  bool empty() {
    return pq.empty();
  void add_lazy(T v) {
     lazy += v;
  }
|};
```

4 Graph

4.1 2-Edge-Connected Components

```
struct EBCC {
   int n, cnt = 0, T = 0;
   vector<vector<int>> adj, comps;
   vector<int>> stk, dfn, low, id;
   EBCC(int n) : n(n), adj(n), dfn(n, -1),
        low(n), id(n, -1) {}
   void addEdge(int u, int v) { adj[u].
        push_back(v), adj[v].push_back(u);
      }
   void build() { for (int i = 0; i < n; i
        ++) { if (dfn[i] == -1) { dfs(i,
        -1); }}}
  void dfs(int u, int p) {
      dfn[u] = low[u] = T++;
   }
}</pre>
```

```
stk.push_back(u);
     for (auto v : adj[u]) {
   if (v == p) { continue; }
        if (dfn[v] == -1) {
          dfs(v, u);
       low[u] = min(low[u], low[v]);
} else if (id[v] == -1) {
          low[u] = min(low[u], dfn[v]);
     if (dfn[u] == low[u]) {
       int x;
       comps.emplace_back();
       do {
          x = stk.back();
          comps.back().push_back(x);
          id[x] = cnt;
          stk.pop_back();
        } while (x != u);
       cnt++;
  }
};
```

4.2 2-Vertex-Connected Components

```
// is articulation point if appear in >=
auto dfs = [&](auto dfs, int u, int p) ->
      void {
   dfn[u] = low[u] = T++;
   for (auto v : adj[u]) {
     if (v == p) { continue; }
if (dfn[v] == -1) {
       stk.push_back(v);
       dfs(dfs, v, u);
       low[u] = min(low[u], low[v]);
       if (low[v] \rightarrow dfn[u]) {
         comps.emplace_back();
         int x;
         do {
           x = stk.back();
           cnt[x]++;
           stk.pop_back();
         } while (x != v);
         comps.back().push_back(u);
         cnt[u]++;
    } else {
       low[u] = min(low[u], dfn[v]);
for (int i = 0; i < n; i++) {
  if (!adj[i].empty()) {
     dfs(dfs, i, -1);
  } else {
     comps.push_back({i});
  }
| }
```

4.3 3-Edge-Connected Components

```
// DSU
struct ETCC {
  int n, cnt = 0;
  vector<vector<int>> adj, comps;
  vector<int> in, out, low, up, nx, id; ETCC(int n): n(n), adj(n), in(n, -1)
       out(in), low(n), up(n), nx(in), id
       (in) {}
  void addEdge(int u, int v) {
    adj[u].push_back(v);
    adj[v].push_back(u);
  void build() {
    int T = 0;
    DSU d(n);
    auto merge = [&](int u, int v) {
      d.join(u, v);
      up[u] += up[v];
    auto dfs = [&](auto dfs, int u, int p
         ) -> void {
```

```
in[u] = low[u] = T++;
     for (auto v : adj[u]) {
  if (v == u) { continue; }
        if (v == p) {
          p = -1;
           continue;
        if (in[v] == -1) {
  dfs(dfs, v, u);
           if (nx[v] == -1 \&\& up[v] <= 1)
             up[u] += up[v];
             low[u] = min(low[u], low[v]);
             continue:
          if (up[v] == 0) { v = nx[v]; }
if (low[u] > low[v]) { low[u] =
           low[v], swap(nx[u], v); } while (v != -1) { merge(u, v);
                 v = nx[v]; }
        } else if (in[v] < in[u]) {</pre>
           low[u] = min(low[u], in[v]);
          up[u]++;
        } else {
           for (int &x = nx[u]; x != -1 &&
                  in[x] \leftarrow in[v] \& in[v] <
                  out[x]; x = nx[x]) {
             merge(u, x);
           up[u]--;
       }
     }
     out[u] = T;
   for (int i = 0; i < n; i++) { if (in[
    i] == -1) { dfs(dfs, i, -1); }}
for (int i = 0; i < n; i++) { if (d.</pre>
         find(i) == i) { id[i] = cnt++;}
  comps.resize(cnt);
   for (int i = 0; i < n; i++) { comps[</pre>
         id[d.find(i)]].push_back(i); }
}
```

4.4 Heavy-Light Decomposition

};

```
struct HLD {
  int n, cur = 0;
  vector<int> sz, top, dep, par, tin,
       tout, seq;
  vector<vector<int>> adj;
  HLD(int n) : n(n), sz(n, 1), top(n),
       dep(n), par(n), tin(n), tout(n),
       seq(n), adj(n) {}
  void addEdge(int u, int v) { adj[u].
       push_back(v), adj[v].push_back(u);
  void build(int root = 0) {
    top[root] = root, dep[root] = 0, par[
        root] = -1;
    dfs1(root), dfs2(root);
  void dfs1(int u) {
    if (auto it = find(adj[u].begin(),
         adj[u].end(), par[u]); it != adj
         [u].end()) {
      adj[u].erase(it);
    for (auto &v : adj[u]) {
      par[v] = u;
      dep[v] = dep[u] + 1;
      dfs1(v);
      sz[u] += sz[v];
      if (sz[v] > sz[adj[u][0]]) { swap(v
           , adj[u][0]); }
   }
  void dfs2(int u) {
    tin[u] = cur++;
    seq[tin[u]] = u;
    for (auto v : adj[u]) {
      top[v] = v == adj[u][0] ? top[u] :
      dfs2(v);
```

```
tout[u] = cur - 1;
  }
  int lca(int u, int v) {
    while (top[u] != top[v]) {
      if (dep[top[u]] > dep[top[v]]) {
        u = par[top[u]];
      } else {
        v = par[top[v]];
      }
    }
    return dep[u] < dep[v] ? u : v;</pre>
   int dist(int u, int v) { return dep[u]
       + dep[v] - 2 * dep[lca(u, v)]; }
  int jump(int u, int k) {
    if (dep[u] < k) { return -1; }</pre>
     int d = dep[u] - k;
     while (dep[top[u]] > d) { u = par[top
         [u]]; }
    return seq[tin[u] - dep[u] + d];
  // u is v's ancestor
  bool isAncestor(int u, int v) { return
       tin[u] <= tin[v] && tin[v] <= tout
       [u]; }
   // root's parent is itself
  int rootedParent(int r, int u) {
     if (r == u) { return u; }
     if (isAncestor(r, u)) { return par[u
         ]; }
     auto it = upper_bound(adj[u].begin(),
          adj[u].end(), r, [\&](int x, int
          y) {
      return tin[x] < tin[y];</pre>
    }) - 1;
return *it;
   // rooted at u, v's subtree size
  int rootedSize(int r, int u) {
     if (r == u) { return n; }
     if (isAncestor(u, r)) { return sz[u];
     return n - sz[rootedParent(r, u)];
  int rootedLca(int r, int a, int b) {
       return lca(a, b) ^ lca(a, r) ^ lca
       (b, r); }
};
```

4.5 Centroid Decomposition

```
vector<int> sz(n), vis(n);
auto build = [&](auto build, int u, int p
    ) -> void {
  sz[u] = 1;
  for (auto v : g[u]) {
    if (v != p && !vis[v]) {
      build(build, v, u);
      sz[u] += sz[v];
 }
};
auto find = [&](auto find, int u, int p,
    int tot) -> int {
  for (auto v : g[u]) {
    if (v != p && !vis[v] && 2 * sz[v] >
         tot) {
      return find(find, v, u, tot);
    }
  return u;
};
auto dfs = [&](auto dfs, int cen) -> void
  build(build, cen, -1);
  cen = find(find, cen, -1, sz[cen]);
  vis[cen] = 1;
  build(build, cen, -1);
  for (auto v : g[cen]) {
    if (!vis[v]) {
      dfs(dfs, v);
```

}

```
|};
|dfs(dfs, 0);
```

4.6 Strongly Connected Components

```
struct SCC {
   int n, cnt = 0, cur = 0;
   vector<int> id, dfn, low, stk;
   vector<vector<int>> adj, comps;
   void addEdge(int u, int v) { adj[u].
        push_back(v); }
  SCC(int n): n(n), id(n, -1), dfn(n, -1), low(n, -1), adj(n) {}
   void build() {
     auto dfs = [&](auto dfs, int u) ->
          void {
       dfn[u] = low[u] = cur++;
       stk.push_back(u);
       for (auto v : adj[u]) {
         if (dfn[v] == -1) {
           dfs(dfs, v);
low[u] = min(low[u], low[v]);
         } else if (id[v] == -1) {
           low[u] = min(low[u], dfn[v]);
       if (dfn[u] == low[u]) {
         int v;
         comps.emplace_back();
         do {
           v = stk.back();
           comps.back().push_back(v);
           id[v] = cnt;
           stk.pop_back();
         } while (u != v);
         cnt++;
       }
     for (int i = 0; i < n; i++) { if (dfn
          [i] == -1) \{ dfs(dfs, i); \} 
     for (int i = 0; i < n; i++) { id[i] =
           cnt - 1 - id[i]; }
     reverse(comps.begin(), comps.end());
   // the comps are in topological sorted
        order
|};
```

4.7 2-SAT

```
struct TwoSat {
 int n, N;
 vector<vector<int>> adj;
  vector<int> ans;
  TwoSat(int n) : n(n), N(n), adj(2 * n)
       {}
 void addClause(int u, bool x) { adj[2 *
 u + !x].push_back(2 * u + x); }
// u == x || v == y
 void addClause(int u, bool x, int v,
    bool y) {
adj[2 * u + !x].push_back(2 * v + y);
    adj[2 * v + !y].push_back(2 * u + x);
 }
// u == x -> v == y
void addImply(int u, bool x, int v,
       bool y) { addClause(u, !x, v, y);
  void addVar() {
    adj.emplace_back(), adj.emplace_back
         ();
 }
 // at most one in var is true
 // adds prefix or as supplementary
       variables
 void atMostOne(const vector<pair<int,</pre>
       bool>> &vars) {
    int sz = vars.size();
    for (int i = 0; i < sz; i++) {
      addVar();
      auto [u, x] = vars[i];
      addImply(u, x, N - 1, true);
      if (i > 0) {
```

```
addImply(N - 2, true, N - 1, true
        );
    addClause(u, !x, N - 2, false);
}

}

// does not return supplementary
    variables from atMostOne()
bool satisfiable() {
    // run tarjan scc on 2 * N
    for (int i = 0; i < 2 * N; i++) { if
        (dfn[i] == -1) { dfs(dfs, i); }}

for (int i = 0; i < N; i++) { if (id
        [2 * i] == id[2 * i + 1]) {
        return false; }}

ans.resize(n);
for (int i = 0; i < n; i++) { ans[i]
        = id[2 * i] > id[2 * i + 1]; }

return true;
}
```

4.8 count 3-cycles and 4cycles

```
sort(ord.begin(), ord.end(), [&](auto i,
       auto j) { return pair(deg[i], i) >
       pair(deg[j], j); });
 for (int i = 0; i < n; i++) { rnk[ord[i]]</pre>
 if (rnk[u] < rnk[v]) { dag[u].push_back(v</pre>
 ); }
// c3
 for (int x = 0; x < n; x++) {
  for (auto y : dag[x]) { vis[y] = 1; }
    for (auto y : dag[x]) { for (auto z :
          dag[y]) { ans += vis[z]; }}
   for (auto y : dag[x]) { vis[y] = 0; }
 // c4
 for (int x = 0; x < n; x++) {
  for (auto y : dag[x]) {    for (auto z :</pre>
          adj[y]) { if (rnk[z] > rnk[x]) {}
   ans += vis[z]+; }}

for (auto y : dag[x]) { for (auto z :

   adj[y]) { if (rnk[z] > rnk[x]) {
          vis[z]--; }}}
1}
```

4.9 Minimum Mean Cycle

create a new vertex S, connect S to all vertices with arbitrary weight (0). Let $f_i(u)$ be the shortest path from S to u with exactly i edges.

```
ans = \min_{f_{n+1}(i)! = \infty} \max_{j=1}^{n} \frac{f_{n+1}(i) - f_{j}(i)}{n+1-j}
```

4.10 Directed Minimum Spanning Tree

```
// DSU with rollback
template <typename Cost>
struct DMST {
 int n;
 vector<int> s, t, lc, rc, h;
vector<Cost> c, tag;
 DMST(int n) : n(n), h(n, -1) {}
 void addEdge(int u, int v, Cost w) {
   int id = s.size();
   s.push_back(u), t.push_back(v), c.
        push_back(w);
   lc.push_back(-1), rc.push_back(-1);
   tag.emplace_back();
   h[v] = merge(h[v], id);
 DSU d(n);
   Cost res{};
   vector<int> vis(n, -1), path(n), q(n)
        , in(n, -1);
   vis[root] = root;
   vector<pair<int, vector<int>>> cycles
   for (auto r = 0; r < n; ++r) {
```

```
auto u = r, b = 0, w = -1;
       while (!~vis[u]) {
         if (!~h[u]) { return {-1, {}}; }
         push(h[u]);
          int e = h[u];
         res += c[e], tag[h[u]] -= c[e];
h[u] = pop(h[u]);
         q[b] = e, path[b++] = u, vis[u] =
         u = d.find(s[e]);
         if (vis[u] == r) {
            int cycle = -1, e = b;
            do {
              w = path[--b];
              cycle = merge(cycle, h[w]);
            } while (d.join(u, w));
           u = d.find(u);
h[u] = cycle, vis[u] = -1;
cycles.emplace_back(u, vector<</pre>
                 int>(q.begin() + b, q.
                 begin() + e));
         }
       for (auto i = 0; i < b; ++i) { in[d
             .find(t[q[i]])] = q[i]; }
     reverse(cycles.begin(), cycles.end())
     for (const auto &[u, comp] : cycles)
       int count = int(comp.size()) - 1;
       d.back(count);
       int ine = in[u];
       for (auto e : comp) { in[d.find(t[e
            ])] = e; }
       in[d.find(t[ine])] = ine;
     vector<int> par;
     par.reserve(n);
     for (auto i : in) { par.push_back(i != -1 ? s[i] : -1); }
     return {res, par};
   void push(int u) {
     c[u] += tag[u];
if (int l = lc[u]; l != -1) { tag[l]
          += tag[u]; }
     if (int r = rc[u]; r != -1) { tag[r]
          += tag[u]; }
     tag[u] = 0;
   int merge(int u, int v) {
     if (u == -1 || v == -1) { return u !=
           -1 ? u : v; }
     push(u);
     push(v);
     if (c[u] > c[v]) { swap(u, v); }
     rc[u] = merge(v, rc[u]);
     swap(lc[u], rc[u]);
     return ū;
   int pop(int u) {
     push(u);
     return merge(lc[u], rc[u]);
};
```

4.11 Maximum Clique

```
pair<int, vector<int>> maxClique(int n,
     const vector<bitset<N>> adj) {
  int mx = 0:
  vector<int> ans, cur;
  auto rec = [&](auto rec, bitset<N> s)
        -> void {
    int sz = s.count();
if (int(cur.size()) > mx) { mx = cur.
          size(), ans = cur; }
     if (int(cur.size()) + sz <= mx) {</pre>
          return; }
    int e1 = -1, e2 = -1;
vector<int> d(n);
     for (int i = 0; i < n; i++) {
       if (s[i]) {
         d[i] = (adj[i] & s).count();
         if (e1 == -1 || d[i] > d[e1]) {
              e1 = i; }
```

```
if (e2 == -1 \mid | d[i] < d[e2]) {
           e2 = i; }
    }
  if (d[e1] >= sz - 2) {
    cur.push_back(e1);
    auto s1 = adj[e1] & s;
    rec(rec, s1);
    cur.pop_back();
    return;
  cur.push_back(e2);
 auto s2 = adj[e2] & s;
rec(rec, s2);
  cur.pop_back();
  s.reset(e2);
  rec(rec, s);
bitset<N> all:
for (int i = 0; i < n; i++) {
 all.set(i);
rec(rec, all);
return pair(mx, ans);
```

4.12 Dominator Tree

```
// res : parent of each vertex in
     dominator tree, -1 is root, -2 if
     not in tree
struct DominatorTree {
  int n, cur = 0;
  vector<int> dfn, rev, fa, sdom, dom,
       val, rp, res;
  vector<vector<int>> adj, rdom, r;
  DominatorTree(int n) : n(n), dfn(n, -1)
         res(n, -2), adj(n), rdom(n), r(n)
       ,
) {
    rev = fa = sdom = dom = val = rp =
         dfn;
  void addEdge(int u, int v) {
    adj[u].push_back(v);
  void dfs(int u) {
    dfn[u] = cur;
    rev[cur] = u;
    fa[cur] = sdom[cur] = val[cur] = cur;
    cur++:
     for (int v : adj[u]) {
      if (dfn[v] == -1) {
        dfs(v);
         rp[dfn[v]] = dfn[u];
      r[dfn[v]].push_back(dfn[u]);
    }
  int find(int u, int c) {
    if (fa[u] == u) { return c != 0 ? -1
         : u; }
    int p = find(fa[u], 1);
    if (p == -1) { return c != 0 ? fa[u]
          : val[u]; }
    if (sdom[val[u]] > sdom[val[fa[u]]])
         { val[u] = val[fa[u]]; }
    fa[u] = p;
    return c != 0 ? p : val[u];
  void build(int s = 0) {
    dfs(s):
     for (int i = cur - 1; i >= 0; i--) {
      for (int u : r[i]) { sdom[i] = min(
           sdom[i], sdom[find(u, 0)]); }
      if (i > 0) { rdom[sdom[i]].
           push_back(i); }
       for (int u : rdom[i]) {
         int p = find(u, 0);
         if (sdom[p] == i) {
           dom[u] = i;
         } else {
           dom[u] = p;
      if (i > 0) { fa[i] = rp[i]; }
```

e[i] = pair(u, v + a), deg[u]++, deg[v +

int col = *max_element(deg.begin(), deg.

vector has(a + b, vector<pair<int, int>>(

4.13 Edge Coloring

// bipartite

a]++;

end());

vector<int> ans(m, -1);

```
col, {-1, -1}));
for (int i = 0; i < m; i++) {
 auto [u, v] = e[i];
 vector<int> c;
  for (auto x : \{u, v\}) {
    c.push_back(0);
    while (has[x][c.back()].first != -1)
         { c.back()++; }
  if (c[0] != c[1]) {
    auto dfs = [\&](auto dfs, int u, int x)
         ) -> void {
      auto [v, i] = has[u][c[x]];
if (v != -1) {
        if (has[v][c[x ^ 1]].first != -1)
          dfs(dfs, v, x \wedge 1);
        } else {
          has[v][c[x]] = \{-1, -1\};
        has[u][c[x ^ 1]] = \{v, i\}, has[v]
             ][c[x \wedge 1]] = \{u, i\};
        ans[i] = c[x \land 1];
      }
    dfs(dfs, v, 0);
 has[u][c[0]] = {v, i};
has[v][c[0]] = {u, i};
 ans[i] = c[0];
// general
auto vizing(int n, const vector<pair<int,</pre>
      int>> &e) {
  vector<int> deg(n);
 for (auto [u, v] : e) {
    deg[u]++, deg[v]++;
  int col = *max_element(deg.begin(), deg
       .end()) + 1;
  vector<int> free(n);
 vector ans(n, vector<int>(n, -1));
  vector at(n, vector<int>(col, -1));
 auto update = [&](int u) {
    free[u] = 0;
    while (at[u][free[u]] != -1) {
      free[u]++;
    }
 };
 auto color = [&](int u, int v, int c1)
    int c2 = ans[u][v];
    ans[u][v] = ans[v][u] = c1;
    at[u][c1] = v, at[v][c1] = u;
    if (c2 != -1) {
      at[u][c2] = at[v][c2] = -1;
      free[u] = free[v] = c2;
    } else {
      update(u), update(v);
    return c2;
 };
 auto flip = [&](int u, int c1, int c2)
    int v = at[u][c1];
    swap(at[u][c1], at[u][c2]);
    if (v != -1) {
      ans[u][v] = ans[v][u] = c2;
```

```
if (at[u][c1] == -1) {
    free[u] = c1;
  if (at[u][c2] == -1) {
    free[u] = c2;
  return v:
for (int i = 0; i < int(e.size()); i++)</pre>
  auto [u, v1] = e[i];
  int v2 = v1, c1 = free[u], c2 = c1, d
  vector<pair<int, int>> fan;
  vector<int> vis(col);
  while (ans[u][v1] == -1) {
    fan.emplace_back(v2, d = free[v2]);
    if (at[v2][c2] == -1) {
      for (int j = int(fan.size()) - 1;
            j >= 0; j--) {
        c2 = color(u, fan[j].first, c2)
    } else if (at[u][d] == -1) {
      for (int j = int(fan.size()) - 1;
            j >= 0; j--) {
        color(u, fan[j].first, fan[j].
             second);
    } else if (vis[d] == 1) {
     break;
    } else {
      vis[d] = 1, v2 = at[u][d];
    }
  if (ans[u][v1] == -1) {
    while (v2 != -1) {
      v2= flip(v2, c2, d);
      swap(c2, d);
    if (at[u][c1] != -1) {
      int j = int(fan.size()) - 2;
      while (j >= 0 && fan[j].second !=
           c2) {
        j--;
      while (j >= 0) {
        color(u, fan[j].first, fan[j].
            second);
      }
   } else {
      i--;
    }
 }
return pair(col, ans);
```

5 String

5.1 Prefix Function

```
template <typename T>
vector<int> prefixFunction(const T &s) {
  int n = int(s.size());
  vector<int> p(n);
  for (int i = 1; i < n; i++) {
    int j = p[i - 1];
    while (j > 0 && s[i] != s[j]) { j = p
        [j - 1]; }
    if (s[i] == s[j]) { j++; }
    p[i] = j;
  }
  return p;
}
```

5.2 Z Function

```
template <typename T>
vector<int> zFunction(const T &s) {
  int n = int(s.size());
  if (n == 0) return {};
  vector<int> z(n);
```

5.3 Suffix Array // need to discretize

```
struct SuffixArray {
 int n;
  vector<int> sa, as, ha;
template <typename T>
 vector<int> sais(const T &s) {
    int n = s.size(), m = *max_element(s.
         begin(), s.end()) + 1;
    vector<int> pos(m + 1), f(n);
    for (auto ch : s) { pos[ch + 1]++; }
    for (int i = 0; i < m; i++) { pos[i +
          1] += pos[i]; }
    for (int i = n - 2; i >= 0; i--) { f[ i] = s[i] != s[i + 1] ? s[i] < s
         [i + 1] : f[i + 1]; 
    vector<int> x(m), sa(n);
    auto induce = [&](const vector<int> &
         ls) {
      fill(sa.begin(), sa.end(), -1);
auto L = [&](int i) { if (i >= 0 &&
            !f[i]) { sa[x[s[i]]++] = i;}
           }};
      auto S = [\&](int i) { if (i >= 0 \&\& )}
            f[i]) { sa[--x[s[i]]] = i;}
           }};
      for (int i = 0; i < m; i++) { x[i]
           = pos[i + 1]; }
      for (int i = int(ls.size()) - 1; i
           >= 0; i--) { S(ls[i]); }
      for (int i = 0; i < m; i++) { x[i]
           = pos[i]; }
      L(n - 1);
      for (int i = 0; i < n; i++) { L(sa[
           i] - 1); }
      for (int i = 0; i < m; i++) { x[i]
      = pos[i + 1]; }
for (int i = n - 1; i >= 0; i--) {
           S(sa[i] - 1); }
    auto ok = [\&](int i) \{ return i == n \}
         | !f[i - 1] && f[i]; };
    auto same = [&](int i, int j) {
      do { if (s[i++] != s[j++]) { return
            false; }} while (!ok(i) && !
           ok(j));
      return ok(i) && ok(j);
    vector<int> val(n), lms;
    for (int i = 1; i < n; i++) { if (ok(</pre>
         i)) { lms.push_back(i); }}
    induce(lms);
    if (!lms.empty()) {
      int p = -1, w = 0;
      for (auto v : sa) {
        if (v != 0 && ok(v)) {
          if (p != -1 \&\& same(p, v)) \{ w \}
                --; }
          val[p = v] = w++;
        }
      auto b = lms;
      for (auto &v : b) { v = val[v]; }
      b = sais(b);
      for (auto &v : b) { v = lms[v]; }
      induce(b);
    return sa;
template <typename T>
 SuffixArray(const T &s) : n(s.size()),
       sa(sais(s)), as(n), ha(n - 1) {
    for (int i = 0; i < n; i++) { as[sa[i
         ]] = i; }
```

```
for (int i = 0, j = 0; i < n; ++i) {
   if (as[i] == 0) {
      j = 0;
   } else {
      for (j -= j > 0; i + j < n && sa[
            as[i] - 1] + j < n && s[i +
            j] == s[sa[as[i] - 1] + j];
            ) { ++j; }
      ha[as[i] - 1] = j;
      }
}
}
};</pre>
```

5.4 Manacher's Algorithm

5.5 Aho-Corasick Automaton

```
constexpr int K = 26;
struct Node {
  array<int, K> nxt;
  int fail = -1;
// other vars
  Node() { nxt.fill(-1); }
vector<Node> aho(1);
for (int i = 0; i < n; i++) {
  string s;
  cin >> s;
  int u = 0;
  for (auto ch : s) {
    int c = ch - 'a'
    if (aho[u].nxt[c] == -1) {
       aho[u].nxt[c] = aho.size();
       aho.emplace_back();
    u = aho[u].nxt[c];
  }
vector<int> q;
for (auto &i : aho[0].nxt) {
  if (i == -1) {
    i = 0;
  } else {
    q.push_back(i);
    aho[i].fail = 0;
for (int i = 0; i < int(q.size()); i++) {</pre>
  int u = q[i];
  if (u > 0) {
    // maintain
  for (int c = 0; c < K; c++) {
    if (int v = aho[u].nxt[c]; v != -1) {
   aho[v].fail = aho[aho[u].fail].nxt[
            c];
       q.push_back(v);
    } else {
       aho[u].nxt[c] = aho[aho[u].fail].
            nxt[c];
```

1}

5.6 Suffix Automaton

```
struct SAM {
  static constexpr int A = 26;
  struct Node {
    int len = 0, link = -1, cnt = 0;
    array<int, A> nxt;
    Node() { nxt.fill(-1); }
   vector<Node> t
  SAM() : t(1) {}
   int size() { return t.size(); }
  Node& operator[](int i) { return t[i];
   int newNode() {
    t.emplace_back();
    return t.size() - 1;
  int extend(int p, int c) {
     int cur = newNode();
     t[cur].len = t[p].len + 1;
     t[cur].cnt = 1;
     while (p != -1 && t[p].nxt[c] == -1)
       t[p].nxt[c] = cur;
      p = t[p].link;
     if (p == -1) {
      t[cur].link = 0;
      else {
       int q = t[p].nxt[c];
      if (t[p].len + 1 == t[q].len) {
  t[cur].link = q;
      } else {
         int clone = newNode();
         t[clone].len = t[p].len + 1;
         t[clone].link = t[q].link;
         t[clone].nxt = t[q].nxt;
         while (p != -1 && t[p].nxt[c] ==
             q) {
           t[p].nxt[c] = clone;
          p = t[p].link;
         t[q].link = t[cur].link = clone;
      }
     return cur;
  }
};
```

5.7 Lexicographically Smallest Rotation

```
template <typename T>
T minRotation(T s) {
  int n = s.size();
  int i = 0, j = 1;
  s.insert(s.end(), s.begin(), s.end());
  while (i < n \& j < n) {
    int k = 0;
    while (k < n \&\& s[i + k] == s[j + k])
    if (s[i + k] \le s[j + k]) {
      j += k + 1;
    } else {
      i += k + 1;
    if (i == j) {
      j++;
  int ans = i < n ? i : j;
  return T(s.begin() + ans, s.begin() +
       ans + n;
```

5.8 EER Tree

```
// cnt : occurrences, (dfs fail tree)
// num : number of pal ending here
struct PAM {
   static constexpr int A = 26;
   struct Node {
    int len = 0, link = 0, cnt = 0, num =
      0;
```

```
array<int, A> nxt{};
     Node() {}
   vector<Node> t;
   int suf = 1;
   string s;
   PAM() : t(2) { t[0].len = -1; }
   int size() { return t.size(); }
   Node& operator[](int i) { return t[i];
   int newNode() {
     t.emplace_back();
     return t.size() - 1;
   bool add(int c, char offset = 'a') {
     int pos = s.size();
     s += c + offset;
    int cur = suf, curlen = 0;
while (true) {
       curlen = t[cur].len;
       if (pos - 1 - curlen >= 0 && s[pos
            -1 - curlen] == s[pos]) {
            break; }
       cur = t[cur].link;
     if (t[cur].nxt[c]) {
       suf = t[cur].nxt[c];
       t[suf].cnt++;
       return false;
     suf = newNode();
     t[suf].len = t[cur].len + 2;
     t[suf].cnt = t[suf].num = 1;
     t[cur].nxt[c] = suf;
     if (t[suf].len == 1) {
       t[suf].link = 1;
return true;
     while (true) {
       cur = t[cur].link;
       curlen = t[cur].len;
       if (pos - 1 - curlen >= 0 && s[pos
            - 1 - curlen] == s[pos]) {
         t[suf].link = t[cur].nxt[c];
         break;
     t[suf].num += t[t[suf].link].num;
     return true;
|};
```

6 Math

6.1 Extended GCD

```
| array<i64, 3> extgcd(i64 a, i64 b) {
| if (b == 0) { return {a, 1, 0}; }
| auto [g, x, y] = extgcd(b, a % b);
| return {g, y, x - a / b * y};
| }
```

6.2 Chinese Remainder Theorem

```
| / /  returns (rem, mod), n = 0 return (0,
     1), no solution return (0, 0)
pair<i64, i64> crt(vector<i64> r, vector<
     i64 > m) {
   int n = r.size();
   for (int i = 0; i < n; i++) {
    r[i] %= m[i];
    if (r[i] < 0) { r[i] += m[i]; }</pre>
   i64 r0 = 0, m0 = 1;
   for (int i = 0; i < n; i++) {
    i64 r1 = r[i], m1 = m[i];
    if (m0 < m1) { swap(r0, r1), swap(m0,</pre>
          m1); }
     if (m0 \% m1 == 0) {
       if (r0 % m1 != r1) { return {0, 0};
       continue;
    auto [g, a, b] = extgcd(m0, m1);
```

```
i64 u1 = m1 / g;
if ((r1 - r0) % g != 0) { return {0,
       0}; }
i64 x = (r1 - r0) / g % u1 * a % u1;
r0 += x * m0;
m0 *= u1;
if (r0 < 0) { r0 += m0; }
}
return {r0, m0};
}</pre>
```

6.3 NTT and polynomials

template <int P>

```
struct Modint {
  int v:
  // need constexpr, constructor, +-*,
       qpow, inv()
template<int P>
constexpr Modint<P> findPrimitiveRoot() {
 Modint<P> i = 2;
int k = __builtin_ctz(P - 1);
  while (true) {
    if (i.qpow((P - 1) / 2).v != 1) {
         break; }
    i = i + 1;
 }
  return i.qpow(P - 1 >> k);
}
template <int P>
constexpr Modint<P> primitiveRoot =
     findPrimitiveRoot<P>();
vector<int> rev;
template <int P>
vector<Modint<P>> roots{0, 1};
template <int P>
void dft(vector<Modint<P>> &a) {
  int n = a.size();
  if (n == 1) { return; }
  if (int(rev.size()) != n) {
    int k = __builtin_ctz(n) - 1;
    rev.resize(n);
    for (int i = 0; i < n; i++) { rev[i]
         = rev[i >> 1] >> 1 | (i & 1) <<
  for (int i = 0; i < n; i++) { if (rev[i
       ] < i) { swap(a[i], a[rev[i]]); }}
  if (roots<P>.size() < n) {</pre>
    int k = __builtin_ctz(roots<P>.size()
    roots<P>.resize(n);
    while ((1 << k) < n) {
      auto e = Modint<P>(primitiveRoot<P</pre>
           >).qpow(P - 1 >> k + 1);
      for (int i = 1 << k - 1; i < 1 << k
        ; i++) {
roots<P>[2 * i] = roots<P>[i];
        roots<P>[2 * i + 1] = roots<P>[i];
* e;
      k++;
    }
  for (int k = 1; k < n; k *= 2) {
  for (int i = 0; i < n; i += 2 * k) {</pre>
      for (int j = 0; j < k; j++) {
        Modint<P> u = a[i + j];
Modint<P> v = a[i + j + k] *
             roots<P>[k + j];
                            j + k] * roots
           fft : v = a[i +
            [n / (2 * k) * j]
        a[i + j] = u + v;
        a[i + j + k] = u - v;
    }
 }
template <int P>
void idft(vector<Modint<P>> &a) {
  int n = a.size();
  reverse(a.begin() + 1, a.end());
  dft(a):
  Modint < P > x = (1 - P) / n;
```

```
for (int i = 0; i < n; i++) { a[i] = a[
       i] * x; }
template <int P>
struct Poly : vector<Modint<P>>> {
 using Mint = Modint<P>;
 Poly() {}
 explicit Poly(int n) : vector<Mint>(n)
  explicit Poly(const vector<Mint> &a) :
      vector<Mint>(a) {}
 explicit Poly(const initializer_list<</pre>
      Mint> &a) : vector<Mint>(a) {}
template<class F>
  explicit Poly(int n, F f) : vector<Mint</pre>
      >(n) { for (int i = 0; i < n; i++) }
        { (*this)[i] = f(i); }}
template<class InputIt>
 explicit constexpr Poly(InputIt first,
      InputIt last) : vector<Mint>(first
        last) {}
 Poly mulxk(int k) {
    auto b = *this;
    b.insert(b.begin(), k, 0);
    return b;
 Poly modxk(int k) {
    k = min(k, int(this->size()));
    return Poly(this->begin(), this->
         begin() + k);
 Poly divxk(int k) {
    if (this->size() <= k) { return Poly</pre>
         (); }
    return Poly(this->begin() + k, this->
         end());
 friend Poly operator+(const Poly &a,
       const Poly &b) {
    Poly res(max(a.size(), b.size()));
    for (int i = 0; i < int(a.size()); i</pre>
         ++) { res[i] = res[i] + a[i]; }
    for (int i = 0; i < int(b.size()); i</pre>
         ++) { res[i] = res[i] + b[i]; }
    return res;
 friend Poly operator-(const Poly &a,
      const Poly &b) {
    Poly res(max(a.size(), b.size()));
    for (int i = 0; i < int(a.size()); i</pre>
         ++) { res[i] = res[i] + a[i]; }
    for (int i = 0; i < int(b.size()); i</pre>
         ++) { res[i] = res[i] - b[i]; }
    return res;
  friend Poly operator*(Poly a, Poly b) {
    if (a.empty() || b.empty()) { return
         Poly(); }
    int sz = 1, tot = a.size() + b.size()
    while (sz < tot) { sz *= 2; }</pre>
    a.resize(sz):
    b.resize(sz);
    dft(a);
    dft(b);
    for (int i = 0; i < sz; i++) { a[i] =
          a[i] * b[i]; }
    idft(a);
    a.resize(tot);
    return a;
 friend Poly operator*(Poly a, Mint b) {
    for (int i = 0; i < int(a.size()); i</pre>
         ++) { a[i] = a[i] * b; }
    return a:
 Poly derivative()
    if (this->empty()) { return Poly(); }
    Poly res(this->size() - 1);
    for (int i = 0; i < this->size() - 1;
    ++i) { res[i] = (i + 1) * (*
         this)[i + 1]; }
    return res;
 Poly integral() {
    Poly res(this->size() + 1);
```

```
for (int i = 0; i < this->size(); ++i
    ) { res[i + 1] = (*this)[i] *
    Mint(i + 1).inv(); }
   return res;
Poly inv(int m) {
   // a[0] != 0
  Poly x({(*this)[0].inv()});
  int k = 1;
  while (k < m) {
   k *= 2;
   x = (x * (Poly({2}) - modxk(k) * x)</pre>
           ).modxk(k);
  return x.modxk(m);
Poly log(int m) {
  return (derivative() * inv(m)).
        integral().modxk(m);
Poly exp(int m) {
  Poly x(\{1\});
   int k = 1;
  while (k < m) {
    k *= 2;
     x = (x^* (Poly(\{1\}) - x.log(k) +
           modxk(k))).modxk(k);
  return x.modxk(m);
Poly pow(i64 k, int m) {
  if (k == 0) { return Poly(m, [&](int
        i) { return i == 0; }); }
   int i = 0;
  while (i < this->size() && (*this)[i
  | .v == 0) { i++; }
| if (i == this->size() || __int128(i) |
| * k >= m) { return Poly(m); }
  Mint v = (*this)[i];
  auto f = divxk(i) * v.inv();
return (f.log(m - i * k) * k).exp(m -
         i * k).mulxk(i * k) * v.qpow(k)
Poly sqrt(int m) {
  // a[0] == 1, otherwise quadratic
        residue?
  Poly x(\{1\});
  int k = 1;
  while (k < m) {
   k *= 2;</pre>
     x = (x + (modxk(k) * x.inv(k)).
           modxk(k)) * ((P + 1) / 2);
  return x.modxk(m);
Poly mulT(Poly b) const {
   if (b.empty()) { return Poly(); }
   int n = b.size();
  reverse(b.begin(), b.end());
return (*this * b).divxk(n - 1);
vector<Mint> evaluate(vector<Mint> x) {
   if (this->empty()) { return vector<</pre>
        Mint>(x.size()); }
   int n = max(x.size(), this->size());
  vector<Poly> q(4 * n);
  vector<Mint> ans(x.size());
  x.resize(n);
  auto build = [&](auto build, int id,
     int l, int r) -> void {
if (r - l == 1) {
        q[id] = Poly(\{1, -x[l].v\});
     } else {
       int m = (l + r) / 2;
build(build, 2 * id, l, m);
build(build, 2 * id + 1, m, r);
q[id] = q[2 * id] * q[2 * id +
             17;
     }
  build(build, 1, 0, n);
  auto work = [&](auto work, int id,
        int l, int r, const Poly &num)
-> void {
     if (r - l == 1) {
```

```
if (l < int(ans.size())) { ans[l]</pre>
               = num[0]; 
      } else {
         work(work, 2 * id + 1, m, r, num.
              mulT(q[2 * id]).modxk(r - m)
      }
    work(work, 1, 0, n, mulT(q[1].inv(n))
    return ans;
  }
};
template <int P>
Poly<P> interpolate(vector<Modint<P>> x,
     vector<Modint<P>> y) {
  // f(xi) = yi
  int n = x.size();
  vector<Poly<P>> p(4 * n), q(4 * n);
  auto dfs1 = [&](auto dfs1, int id, int
       1, int r) -> void {
    if (l == r) {
      p[id] = Poly < P > ({-x[1].v, 1});
    int m = l + r >> 1;
dfs1(dfs1, id << 1, l, m);
dfs1(dfs1, id << 1 | 1, m + 1, r);
    p[id] = p[id << 1] * p[id << 1 | 1];
  dfs1(dfs1, 1, 0, n - 1);
Poly<P> f = Poly<P>(p[1].derivative().
       evaluate(x));
  auto dfs2 = [&](auto dfs2, int id, int
    l, int r) -> void {
if (l == r) {
      q[id] = Poly < P > ({y[l] * f[l].inv()}
           });
      return;
    int m = 1 + r >> 1;
    dfs2(dfs2, 1, 0, n - 1);
  return q[1];
auto shift = [&](FPS f, int k) {
  FPS a(n + 1), b(n + 1);
  Mint powk = 1;
  for (int i = 0; i <= n; i++) {
    a[i] = ifact[i] * powk;
    b[i] = fact[i] * f[i];
powk = powk * k;
  reverse(b.begin(), b.end());
auto g = a * b;
  g.resize(n + 1);
  reverse(g.begin(), g.end());
  for (int i = 0; i <= n; i++) {
  g[i] = g[i] * ifact[i];</pre>
  return g;
```

6.4 Any Mod NTT

```
constexpr int P0 = 998244353, P1 =
    1004535809, P2 = 469762049;
constexpr i64 P01 = 1LL * P0 * P1;
constexpr int inv0 = Modint<P1>(P0).inv()
    .v;
constexpr int inv01 = Modint<P2>(P01).inv
    ().v;
for (int i = 0; i < int(c.size()); i++) {
    i64 x = 1LL * (c1[i] - c0[i] + P1) % P1
        * inv0 % P1 * P0 + c0[i];
    c[i] = ((c2[i] - x % P2 + P2) % P2 *
        inv01 % P2 * (P01 % P) % P + x) %
    P;
}</pre>
```

6.5 Newton's Method

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

6.6 Fast Walsh-Hadamard Transform

```
1. XOR Convolution
```

- $f(A) = (f(A_0) + f(A_1), f(A_0) f(A_1))$
- $f^{-1}(A) = (f^{-1}(\frac{A_0 + A_1}{2}), f^{-1}(\frac{A_0 A_1}{2}))$
- 2. OR Convolution
 - $f(A) = (f(A_0), f(A_0) + f(A_1))$ • $f^{-1}(A) = (f^{-1}(A_0), f^{-1}(A_1) - f^{-1}(A_0))$
- 3. AND Convolution
 - $f(A) = (f(A_0) + f(A_1), f(A_1))$ • $f^{-1}(A) = (f^{-1}(A_0))$ • $f^{-1}(A_1), f^{-1}(A_1))$

6.7 Simplex Algorithm

Description: maximize $\mathbf{c}^T \mathbf{x}$ subject to $A\mathbf{x} \leq \mathbf{b}$ and $\mathbf{x} \geq 0$. Returns $-\infty$ if infeasible and ∞ if unbounded.

```
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
vector<vector<double>> d;
vector<int> p, q;
void pivot(int r, int s) {
  double inv = 1.0 / d[r][s];
  for (int i = 0; i < m + 2; ++i) {
  for (int j = 0; j < n + 2; ++j) {
    if (i != r && j != s) d[i][j] -= d[
        r][j] * d[i][s] * inv;</pre>
     }
  for (int i = 0; i < m + 2; ++i) if (i
  != r) d[i][s] *= -inv;
for (int j = 0; j < n + 2; ++j) if (j
         != s) d[r][j] *= +inv;
   d[r][s] = inv;
  swap(p[r], q[s]);
bool phase(int z) {
  int x = m + z;
  while (true) {
     int s = -1;
     for (int i = 0; i <= n; ++i) {
        if (!z && q[i] == -1) continue;
if (s == -1 || d[x][i] < d[x][s]) s
     if (d[x][s] > -eps) return true;
     int r = -1;
for (int i = 0; i < m; ++i) {
        if (d[i][s] < eps) continue;
if (r == -1 || d[i][n + 1] / d[i][s</pre>
              ] < d[r][n + 1] / d[r][s]) r =
               i;
     if (r == -1) return false;
     pivot(r, s);
  }
}
vector<double> solve(const vector<vector<</pre>
      double>> &a, const vector<double> &b
       , const vector<double> &c) {
  m = b.size(), n = c.size();
  d = vector<vector<double>>(m + 2,
         vector<double>(n + 2));
  for (int i = 0; i < m; ++i) {
  for (int j = 0; j < n; ++j) d[i][j] =</pre>
             a[i][j];
  p.resize(m), q.resize(n + 1);
for (int i = 0; i < m; ++i) p[i] = n +</pre>
         i, d[i][n] = -1, d[i][n + 1] = b[i]
```

for (int i = 0; i < n; ++i) q[i] = i, d

[m][i] = -c[i];

q[n] = -1, d[m + 1][n] = 1;

```
int r = 0:
for (int i = 1; i < m; ++i) if (d[i][n
     + 1] < d[r][n + 1]) r = i;
if (d[r][n + 1] < -eps) {
  pivot(r, n);
  if (!phase(1) || d[m + 1][n + 1] < -
       eps) return vector<double>(n, -
       inf);
  for (int i = 0; i < m; ++i) if (p[i]
       == -1) {
    int s = min_element(d[i].begin(), d
        [i].end() - 1) - d[i].begin();
    pivot(i, s);
if (!phase(0)) return vector<double>(n,
      inf);
vector<double> x(n);
for (int i = 0; i < m; ++i) if (p[i] <</pre>
     n) x[p[i]] = d[i][n + 1];
```

6.7.1 Construction

Standard form: maximize $\mathbf{c}^T\mathbf{x}$ subject to $A\mathbf{x} \leq \mathbf{b}$ and $\mathbf{x} \geq 0$. Dual LP: minimize $\mathbf{b}^T\mathbf{y}$ subject to $A^T\mathbf{y} \geq \mathbf{c}$ and $\mathbf{y} \geq \mathbf{0}$. $\bar{\mathbf{x}}$ and $\bar{\mathbf{y}}$ are optimal if and only if for all $i \in [1, n]$, either $\bar{x}_i = 0$ or $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$ holds and for all $i \in [1, m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^n A_{ij}\bar{x}_j = b_j$ holds.

```
1. In case of minimization, let c_i' = -c_i

2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le b_j
3. \sum_{1 \le i \le n} A_{ji} x_i = b_j
```

• $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$ • $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$ 4. If x_i has no lower bound, replace x_i with

6.8 Subset Convolution

Description: $h(s) = \sum_{s' \subset s} f(s')g(s \setminus s')$

```
vector<int> SubsetConv(int n, const
       vector<int> &f, const vector<int> &g
       ) {
    const int m = 1 \ll n;
    vector < vector < int>> a(n + 1, vector < int)
          >(m)), b(n + 1, vector<int>(m));
    for (int i = 0; i < m; ++i) {
      a[__builtin_popcount(i)][i] = f[i];
      b[__builtin_popcount(i)][i] = g[i];
    for (int i = 0; i <= n; ++i) {
      for (int j = 0; j < n; ++j) {
  for (int s = 0; s < m; ++s) {
    if (s >> j & 1) {
      a[i][s] += a[i][s ^ (1 << j)];
    }
               b[i][s] += b[i][s \land (1 << j)];
         }
      }
   }
    vector<vector<int>> c(n + 1, vector<int
          >(m));
   for (int s = 0; s < m; ++s) {
  for (int i = 0; i <= n; ++i) {
    for (int j = 0; j <= i; ++j) c[i][s
      ] += a[j][s] * b[i - j][s];
      }
    for (int i = 0; i \le n; ++i) {
      for (int j = 0; j < n; ++j) {
  for (int s = 0; s < m; ++s) {
    if (s >> j & 1) c[i][s] -= c[i][s
                    ^ (1 << j)];
      }
    vector<int> res(m);
   for (int i = 0; i < m; ++i) res[i] = c[
           __builtin_popcount(i)][i];
    return res;
1}
```

6.9 Berlekamp Massey Algorithm

```
// find \sum a_(i-j)c_j = 0 for d <= i template <typename T>
vector<T> berlekampMassey(const vector<T>
       &a) {
  vector<T> c(1, 1), oldC(1);
  int oldI = -1;
T oldD = 1;
  for (int i = 0; i < int(a.size()); i++)</pre>
     T d = 0;
     for (int j = 0; j < int(c.size()); j
++) { d += c[j] * a[i - j]; }
     if (d == 0) { continue; }
T mul = d / oldD;
     vector<T> nc = c;
     nc.resize(max(int(c.size()), i - oldI
     + int(oldC.size())));
for (int j = 0; j < int(oldC.size());
     j++) { nc[j + i - oldI] -= oldC
[j] * mul; }
if (i - int(c.size()) > oldI - int(
           oldC.size())) {
        oldI = i;
        oldD = d
        swap(oldC, c);
     swap(c, nc);
  return c;
```

6.10 Fast Linear Recurrence

```
// p : a[0] ~ a[d - 1]
// q : a[i] = \sum_{a[i - j]q[j]}
template <typename T>
T linearRecurrence(vector<T> p, vector<T>
      q, i64 n) {
  int d = q.size() - 1;
  assert(int(p.size()) == d);
p = p * q;
  p.resize(d):
  while (n > 0) {
    auto nq = q;
    for (int i = 1; i <= d; i += 2) {
      nq[i] *= -1;
    auto np = p * nq;
nq = q * nq;
    for (int i = 0; i < d; i++) {
p[i] = np[i * 2 + n % 2];
    for (int i = 0; i <= d; i++) {
    q[i] = nq[i * 2];
    n /= 2;
  return p[0] / q[0];
```

6.11 Prime check and factorize

```
i64 mul(i64 a, i64 b, i64 mod) {}
i64 qpow(i64 x, i64 p, i64 mod) {}
bool isPrime(i64 n) {
  if (n == 1) { return false; }
  int r = __builtin_ctzll(n - 1);
  i64 d = n - 1 >> r;
  auto checkComposite = [&](i64 p) {
    i64 x = qpow(p, d, n);
    if (x == 1 \mid | x == n - 1) { return
         false; }
    for (int i = 1; i < r; i++) {
      x = mul(x, x, n);
      if (x == n - 1) { return false; }
    return true;
  for (auto p : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}) {
    if (n == p) {
  return true;
    } else if (checkComposite(p)) {
```

```
return false:
    }
  return true;
vector<i64> pollardRho(i64 n) {
  vector<i64> res;
  auto work = [&](auto work, i64 n) {
  if (n <= 10000) {</pre>
       for (int i = 2; i * i <= n; i++) {</pre>
         while (n % i == 0) {
           res.push_back(i);
           n /= i;
       if (n > 1) { res.push_back(n); }
    return;
} else if (isPrime(n)) {
      res.push_back(n);
      return;
    i64 x0 = 2;
auto f = [&](i64 x) { return (mul(x,
         x, n) + 1) % n; };
    while (true) {
      64 \times = x0, y = x0, d = 1, power = 1, lam = 0, v = 1;
      while (d == 1) {
         y = f(y);
         ++lam;
         v = mul(v, abs(x - y), n);
         if (lam % 127 == 0) {
           d = gcd(v, n);
v = 1;
         if (power == lam) {
           x = y;
power *= 2;
           lam = 0;
           d = gcd(v, n);
        }
      if (d != n) {
         work(work, d);
         work(work, n / d);
         return;
       ++x0;
    }
  };
  work(work, n);
  sort(res.begin(), res.end());
  return res;
```

6.12 Count Primes leq n

```
// __attribute__((target("avx2")
      optimize("03", "unroll-loops")))
i64 primeCount(const i64 n) {
  if (n <= 1) { return 0; }
if (n == 2) { return 1; }</pre>
  const int v = sqrtl(n);
  int s = (v + 1) / 2;
  vector<int> smalls(s), roughs(s), skip(
  v + 1);
vector<i64> larges(s);
  iota(smalls.begin(), smalls.end(), 0);
  for (int i = 0; i < s; i++) {
  roughs[i] = 2 * i + 1;</pre>
     larges[i] = (n / roughs[i] - 1) / 2;
  const auto half = [](int n) -> int {
        return (n - 1) >> 1; };
  int pc = 0;
  if (pt p = 3; p <= v; p += 2) {
   if (skip[p]) { continue; }
   int q = p * p;
   if (1LL * q * q > n) { break; }
     skip[p] = true;
     for (int i = q; i \ll v; i += 2 * p)
           skip[i] = true;
     int ns = 0;
     for (int k = 0; k < s; k++) {
       int i = roughs[k];
       if (skip[i]) { continue; }
i64 d = 1LL * i * p;
```

```
larges[ns] = larges[k] - (d <= v ?
    larges[smalls[d / 2] - pc] :
    smalls[half(n / d)]) + pc;</pre>
      roughs[ns++] = i;
  s = ns;
  for (int i = half(v), j = v / p - 1 |
1; j >= p; j -= 2) {
int c = smalls[j / 2] - pc;
for (int e = j * p / 2; i >= e; i
             --) { smalls[i] -= c; }
  pc++;
larges[0] += 1LL * (s + 2 * (pc - 1)) *
(s - 1) / 2;
for (int k = 1; k < s; k++) { larges[0]
-= larges[k]; }
for (int l = 1; l < s; l++) {
  i64 q = roughs[1];
  i64 M = n / q;
int e = smalls[half(M / q)] - pc;
   if (e <= 1) { break; }</pre>
  i64 t = 0;
  for (int k = l + 1; k \le e; k++) { t
         += smalls[half(M / roughs[k])];
  larges[0] += t - 1LL * (e - l) * (pc
         + 1 - 1);
return larges[0] + 1;
```

6.13 Discrete Logarithm

```
| //  return min x >= 0 s.t. a ^ x = b mod m
       , 0 \land 0 = 1, -1 if no solution
 // (I think) if you want x > 0 (m != 1),
      remove if (b == k) return add;
 int discreteLog(int a, int b, int m) {
   if (m == 1) {
     return 0;
   a %= m, b %= m;
   int k = 1, add = 0, g;
   while ((g = gcd(a, m)) > 1) {
     if (b == k) {
        return add;
     } else if (b % g) {
  return -1;
     b /= g, m /= g, ++add;
k = 1LL * k * a / g % m;
   if (b == k) {
     return add;
   int n = sqrt(m) + 1;
   int an = 1;
   for (int i = 0; i < n; ++i) {
    an = 1LL * an * a % m;
   unordered_map<int, int> vals;
   for (int q = 0, cur = b; q < n; ++q) {
     vals[cur] = q;
cur = 1LL * a * cur % m;
   for (int p = 1, cur = k; p <= n; ++p) {
   cur = 1LL * cur * an % m;</pre>
     if (vals.count(cur)) {
        int ans = n * p - vals[cur] + add;
        return ans;
     }
   return -1:
```

6.14 Quadratic Residue

```
// rng
int jacobi(int a, int m) {
  int s = 1;
  while (m > 1) {
    a %= m;
    if (a == 0) { return 0; }
    int r = __builtin_ctz(a);
    if (r % 2 == 1 && (m + 2 & 4) != 0) {
        s = -s; }
```

```
if ((a \& m \& 2) != 0) \{ s = -s; \}
    swap(a, m);
  return s;
int quadraticResidue(int a, int p) {
  if (p == 2) { return a % 2; }
  int j = jacobi(a, p);
  if (j == 0 \mid | j == -1) \{ return j; \}
  int b, d;
  while (true) {
   b = rng() % p;
d = (1LL * b * b + p - a) % p;
    if (jacobi(d, p) == -1) \{ break; \}
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp
  for (int e = p + 1 >> 1; e > 0; e >>=
      1) {
    if (e % 2 == 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * g1
            % p * f1 % p) % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0
           ) % p;
      g0 = tmp;
    tmp = (1LL * f0 * f0 + 1LL * d * f1 %
          p * f1 % p) % p;
    f1 = 2LL * f0 * f1 % p;
    f0 = tmp;
  return g0;
```

6.15 Characteristic Polynomial

```
vector<vector<int>>> Hessenberg(const
     vector<vector<int>> &A) {
  int N = A.size();
  vector<vector<int>> H = A;
  for (int i = 0; i < N - 2; ++i) {
    if (!H[i + 1][i]) {
       for (int j = i + 2; j < N; ++j) {
         if (H[j][i]) {
           for (int k = i; k < N; ++k)
                swap(H[i + 1][k], H[j][k])
           for (int k = 0; k < N; ++k)
                 swap(H[k][i + 1], H[k][j])
           break;
         }
      }
    if (!H[i + 1][i]) continue;
    int val = fpow(H[i + 1][i], kP - 2);
    for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP</pre>
       for (int k = i; k < N; ++k) H[j][k]
= (H[j][k] + 1LL * H[i + 1][k
            ] * (kP - coef)) % kP;
       for (int k = 0; k < N; ++k) H[k][i
            + 1] = (H[k][i + 1] + 1LL * H[
            k][j] * coef) % kP;
    }
  return H;
vector<int> CharacteristicPoly(const
     vector<vector<int>> &A) {
  int N = A.size();
  auto H = Hessenberg(A);
  for (int i = 0; i < N; ++i) {
  for (int j = 0; j < N; ++j) H[i][j] =</pre>
           kP - H[i][j];
  vector<vector<int>>> P(N + 1, vector<int</pre>
       >(N + 1));
  P[0][0] = 1;
  for (int i = 1; i <= N; ++i) {
    P[i][0] = 0;
    for (int j = 1; j <= i; ++j) P[i][j]</pre>
         = P[i - 1][j - 1];
    int val = 1;
```

6.16 Linear Sieve Related

```
vector<int> minp(N + 1), primes, mobius(N
       + 1);
mobius[1] = 1;
for (int i = 2; i \le N; i++) {
  if (!minp[i]) {
    primes.push_back(i);
     minp[i] = i;
    mobius[i] = -1;
   for (int p : primes) {
    if (p > N / i) {
       break:
     minp[p * i] = p;
    mobius[p * i] = -mobius[i];
if (i % p == 0) {
       mobius[p * i] = 0;
       break;
  }
| }
```

6.17 De Bruijn Sequence

```
int res[kN], aux[kN], a[kN], sz;
void Rec(int t, int p, int n, int k) {
  if (t > n) {
    if (n \% p == 0)
       for (int i = 1; i <= p; ++i) res[sz
             ++] = aux[i];
    aux[t] = aux[t - p];
    Rec(t + 1, p, n, k);
for (aux[t] = aux[t - p] + 1; aux[t]
          < k; ++aux[t]) Rec(t + 1, t, n,
 }
int DeBruijn(int k, int n) {
   // return cyclic string of length k^n
        such that every string of length n
         using k character appears as a
        substring.
  if (k == 1) return res[0] = 0, 1;
fill(aux, aux + k * n, 0);
  return sz = 0, Rec(1, 1, n, k), sz;
```

6.18 Floor Sum

6.19More Floor Sum

•
$$m = \lfloor \frac{an+b}{c} \rfloor$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), \\ 0, \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), \end{cases} \end{split}$$

$$h(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^{2}$$

$$= \begin{cases} \lfloor \frac{a}{c} \rfloor^{2} \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^{2} \cdot (n+1) & \frac{cank(D)}{2} \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) & \frac{cank(D)}{2} \end{cases}$$

$$= \begin{cases} -\frac{a}{c} \rfloor^{2} \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^{2} \cdot (n+1) & \frac{cank(D)}{2} \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) & \frac{cank(D)}{2} \end{cases}$$

$$+ 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) & \text{Cayley's} \end{cases}$$

$$+ 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq e \lor \oplus_{i} \\ 0, & n < 0 \lor d_{\mathbf{q}} \\ nm(m+1) - 2g(c, c-b-1, a, m-1) & \text{tic} \\ -2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases}$$

6.20 Min Mod Linear

```
// \min i : [0, n) (a * i + b) % m
// ok in 1e9
 int minModLinear(int n, int m, int a, int
        b, int cnt = 1, int p = 1, int q =
      1) {
   if (a == 0) { return b; } if (cnt % 2 == 1) {
      if (b >= a) {
        int t = (m - b + a - 1) / a;
int c = (t - 1) * p + q;
if (n <= c) { return b; }
n -= c;
b += a * t - m;
     b = a - 1 - b;
   } else {
      if (b < m - a) {
        int t = (m - b - 1) / a;
        b = m - 1 - b;
   }
   cnt++;
   int d = m / a;
   int c = minModLinear(n, a, m % a, b,
   cnt, (d - 1) * p + q, d * p + q);
return cnt % 2 == 1 ? m - 1 - c : a - 1
          - c;
į }
```

Count of subsets with 6.21sum (mod P) leq T

```
int n, T;
cin >> n >> T;
vector<int> cnt(T + 1);
for (int i = 0; i < n; i++) {
  int a;
  cin >> a;
  cnt[a]++;
vector<Mint> inv(T + 1);
for (int i = 1; i <= T; i++) {
  inv[i] = i == 1 ? 1 : -P / i * inv[P %
      i];
FPS f(T + 1);
f = f.exp(T + 1);
```

6.22 Theorem

· Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge

- The number of undirected spanning $a \geq c \vee ih \mathscr{E} ds |\det(\tilde{L}_{11})|.$

 $n < \underline{0} \vee_{\text{The number of directed spanning}} 0$ tree rooted at r in G is $|\det(L_{rr})|$. otherwise

• Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ $(x_{ij} \text{ is chosen uniformly at random})$ if i < jand $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

· Cayley's Formula

 $a \ge \epsilon \lor \text{diven}$ a degree sequence $n < 0 \ \forall d_0, = d_2, \ldots, d_n$ for each labeled vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1, 2, \dots, k$ belong to different components. Then $T \cdot -kn^{n-k-1}$ $T_{n,k} = kn^{n-}$
- Erdős–Gallai Theorem

A sequence of non-negative integers $d_1 \ge$ $d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + d_2 + \ldots + d_n$

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all $1 \leq k \leq n$.

• Burnside's Lemma

Let X be a set and G be a group that acts on X. For $g \in G$, denote by X^g the elements fixed by g:

$$X^g = \{ x \in X \mid gx \in X \}$$

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

• Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and b_1, \ldots, b_n is bigraphic

$$a_1 \ge \cdots \ge a_n$$
 and b_1, \ldots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \le a_i$

 $\sum_{i=1}^{n} \min(b_i, k) \text{ holds for every } 1 \leq k \leq n.$

Sequences a and b called bigraphic if there is a labeled simple bipartite graph such that a and b is the degree sequence of this bipartite graph.

• Fulkerson-Chen-Anstee theorem

A sequence $(a_1, b_1), \ldots, (a_n, b_n)$ of nonnegative integer pairs with $a_1 \geq \cdots \geq$ a_n is digraphic if and only if $\sum a_i =$

$$\sum_{i=1}^{n} b_i \text{ and } \sum_{i=1}^{k} a_i \leq \sum_{i=1}^{k} \min(b_i, k-1) + \sum_{i=k+1}^{n} \min(b_i, k) \text{ holds for every } 1 \leq k \leq n.$$

Sequences a and b called digraphic if there is a labeled simple directed graph such that each vertex v_i has indegree a_i and outdegree b_i .

• Pick's theorem

For simple polygon, when points we have are all integer, we have #{lattice points in the interior} $\#\{\text{lattice points on the boundary}\}\ -1$

• Möbius inversion formula

$$\begin{array}{lll} - \ f(n) &=& \sum_{d\mid n} g(d) & \Leftrightarrow & g(n) &= \\ & \sum_{d\mid n} \mu(d) f(\frac{n}{d}) \\ - \ f(n) &=& \sum_{n\mid d} g(d) & \Leftrightarrow & g(n) &= \\ & \sum_{n\mid d} \mu(\frac{d}{n}) f(d) & & & \end{array}$$

- · Spherical cap
 - A portion of a sphere cut off by a
 - $\begin{array}{ll} \text{plane.} \\ r\text{: sphere radius, } a\text{: radius of the} \\ \text{base of the cap, } h\text{: height of the cap,} \end{array}$

 - base of the cap, it. logic of the cap, the cap,
- The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200000 for n < 1e19.
- The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30 for $n=0\sim 9,$ 627 for $n=20,\sim 2e5$ for $n = 50, \sim 2e8 \text{ for } n = 100.$
- Total number of partitions of n distinct elements: B(n) = 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, 115975, 6785• Total number of 27644437, 190899322,
- Ordinary Generating Function A(x) = $\sum_{i\geq 0} a_i x^i$

$$\begin{array}{l} -A(rx) \Rightarrow r^n a_n \\ -A(x) + B(x) \Rightarrow a_n + b_n \\ -A(x) B(x) \Rightarrow \sum_{i=0}^n a_i b_{n-i} \\ -A(x)^k \Rightarrow \sum_{i_1 + i_2 + \dots + i_k = n} a_{i_1} a_{i_2} \dots a_{i_k} \\ -x A(x)' \Rightarrow n a_n \\ -\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^n a_i \end{array}$$

• Exponential Generating Function A(x) = $\sum_{i>0} \frac{a_i}{i!} x_i$

$$\begin{array}{l} -A(x)+B(x)\Rightarrow a_n+b_n\\ -A^{(k)}(x)\Rightarrow a_{n+k}\\ -A(x)B(x)\Rightarrow \sum_{i=0}^{k_n}\binom{n}{i}a_ib_{n-i}\\ -A(x)^k\Rightarrow \sum_{i_1+i_2+\cdots+i_k=n}\binom{n}{i_1,i_2,\ldots,i_k}a_{i_1}\\ -xA(x)\Rightarrow na_n \end{array}$$

• Special Generating Function

$$- (1+x)^{n} = \sum_{i \ge 0} {n \choose i} x^{i}$$

$$- \frac{1}{(1-x)^{n}} = \sum_{i \ge 0} {n \choose i} x^{i}$$

• Bernoulli numbers

$$B_0 = 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, \text{ EGF is } B(x) = \frac{x}{e^{x}-1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k^+ n^{m+1-k}$$

· Stirling numbers of the second kind Partitions of n distinct elements into exactly k

$$\begin{array}{lll} S(n,k) &=& S(n-1,k-1) \, + \, kS(n-1,k), S(n,1) = S(n,n) = 1 \\ S(n,k) &=& \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n &=& \sum_{i=0}^n S(n,i)(x)_i \end{array}$$

• Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1 - x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \geq j$, k j:s s.t. $\pi(j) > j$. E(n,k) = (n-k)E(n-1,k-1) + (k+1)1)E(n-1,k)E(n,0) = E(n, n-1) = 1

$$E(n, 0) = E(n, n - 1) = 1$$

$$E(n, k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

7 Dynamic Programming

7.1 Dynamic Convex Hull

```
struct Line {
   // kx + b
   mutable i64 k, b, p;
   bool operator<(const Line& o) const {</pre>
        return k < o.k; }</pre>
   bool operator<(i64 x) const { return p</pre>
        < x;  }
 struct DynamicConvexHullMax : multiset<</pre>
      Line, less<>>> {
   // (for doubles, use INF = 1/.0, div(a,
        b) = a/b
   static constexpr i64 INF =
        numeric_limits<i64>::max();
   i64 div(i64 a, i64 b) {
     // floor
     return a / b - ((a \land b) < 0 \&\& a \% b)
   bool isect(iterator x, iterator y) {
     if (y == end()) return x -> p = INF, 0;
     if (x->k == y->k) x->p = x->b > y->b
? INF : -INF;
     else x->p = div(y->b - x->b, x->k - y
          ->k);
     return x->p >= y->p;
   void add(i64 k, i64 b) {
     auto z = insert(\{k, b, 0\}), y = z++,
          X = V;
     while (isect(y, z)) z = erase(z);
if (x != begin() && isect(--x, y))
          isect(x, y = erase(y));
     while ((y = x) != begin() \&\& (--x)->p
           >= y->p)
       isect(x, erase(y));
   i64 query(i64 x) {
     if (empty()) {
       return - INF;
     auto l = *lower_bound(x);
     return 1.k * x + 1.b;
| };
```

7.2 1D/1D Convex Optimization

```
struct segment {
  int i, l, r;
  segment(int a, int b, int c): i(a), l(b
       ), r(c) {}
inline long long f(int l, int r) { return
      dp[l] + w(l + 1, r); }
void solve() {
  dp[0] = 011;
  deque<segment> deq; deq.push_back(
  segment(0, 1, n));
for (int i = 1; i <= n; ++i) {</pre>
  dp[i] = f(deq.front().i, i);
  while (deq.size() && deq.front().r < i</pre>
       + 1) deq.pop_front();
  deq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
  while (deq.size() && f(i, deq.back().1)
        < f(deq.back().i, deq.back().l))
       deq.pop_back();
  if (deq.size()) {
    int d = 1048576, c = deq.back().1;
    while (d \gg 1) if (c + d \ll deq.back)
         ().r) {
    if (f(i, c + d) > f(deq.back().i, c +
          d)) c += d;
    deq.back().r = c; seq.l = c + 1;
  if (seg.l <= n) deq.push_back(seg);</pre>
```

7.3 Condition

7.3.1 Totally Monotone (Con-|} cave/Convex)

```
\begin{array}{ll} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \implies B[i][j'] \leq \\ B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \implies B[i][j'] \geq \\ B[i'][j'] \end{array}
```

7.3.2 Monge Condition (Concave/ Convex)

```
\begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

7.3.3 Optimal Split Point

```
If B[i][j] + B[i+1][j+1] \ge B[i][j+1] + B[i+1][j] then H_{i,j-1} \le H_{i,j} \le H_{i+1,j}
```

8 Ckisemetry

8.1 Basic #define IM imag

```
#define RE real
using lld = int64_t;
using llf = long double;
using PT = complex<lld>;
using PF = complex<llf>;
using P = PT;
llf abs(P p) { return sqrtl(norm(p)); }
PF toPF(PT p) { return PF{RE(p), IM(p)};
int sgn(lld x) { return (x > 0) - (x < 0)
lld dot(P a, P b) { return RE(conj(a) * b
    ); }
lld cross(P a, P b) { return IM(conj(a) *
b); }
int ori(P a, P b, P c) {
 return sgn(cross(b - a, c - a));
int quad(P p) {
 return (IM(p) == 0) // use sgn for PF
   ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ?
        0:2);
int argCmp(P a, P b) {
  // returns 0/+-1, starts from theta = -
  int qa = quad(a), qb = quad(b);
 if (qa != qb) return sgn(qa - qb);
 return sgn(cross(b, a));
P rot90(P p) { return P{-IM(p), RE(p)}; }
template <typename V> llf area(const V &
    pt) {
 11d ret = 0; // BE CAREFUL OF TYPE!
  for (int i = 1; i + 1 < (int)pt.size();</pre>
   ret += cross(pt[i] - pt[0], pt[i+1] -
 pt[0]);
return ret / 2.0;
template <typename V> PF center(const V &
 for (int i = 1; i + 1 < (int)pt.size();</pre>
       i++) {
   ret += (pt[i] + pt[i + 1] + pt[0]) * cur; A += cur;
```

```
return toPF(ret) / llf(A * 3);
}
PF project(PF p, PF q) { // p onto q
return dot(p, q) * q / dot(q, q); //
dot<llf>
}
```

8.2 ConvexHull

```
// from NaCl, counterclockwise, be
    careful of n<=2
vector<P> convex_hull(vector<P> v) { // n
     ==0 will RE
  sort(all(v)); // by X then Y
 if (v[0] == v.back()) return \{v[0]\};
 int t = 0, s = 1; vector<P> h(v.size()
      + 1);
  for (int
            _ = 2; _--; s = t--, reverse(
      all(v)))
    for (P p : v) {
     while (t>s && ori(p, h[t-1], h[t
           -2]) >= 0) t--;
     h[t++] = p;
 return h.resize(t), h;
```

8.3 CyclicTS

8.4 Delaunay

```
/* please ensure input points are unique
/* A triangulation such that no points
    will strictly
inside circumcircle of any triangle. C
    should be big
enough s.t. the initial triangle contains
     all points */
#define L(i) ((i)==0 ? 2 : (i)-1)
#define R(i) ((i)==2 ? 0 : (i)+1)
#define F3 for (int i = 0; i < 3; i++)
bool is_inf(P z) { return RE(z) <= -C ||</pre>
    RE(z) >= C; 
bool in_cc(const array<P,3> &p, P q) {
  i128 inf_det = 0, det = 0, inf_N, N;
  F3 {
    if (is_inf(p[i]) && is_inf(q))
         continue;
    else if (is_inf(p[i])) inf_N = 1, N =
          -norm(q);
    else if (is_inf(q)) inf_N = -1, N =
        norm(p[i]);
    else inf_N = 0, N = norm(p[i]) - norm
        (q);
    lld D = cross(p[R(i)] - q, p[L(i)] -
        q);
    inf_det += inf_N * D; det += N * D;
  return inf_det != 0 ? inf_det > 0 : det
       > 0;
P v[maxn];
struct Tri;
struct E {
  Tri *t; int side;
  E(Tri *t_=0, int side_=0) : t(t_), side
      (side_) {}
struct Tri {
  array<int,3> p; array<Tri*,3> ch; array
       <E,3> e;
```

```
Tri(int a=0, int b=0, int c=0) : p{a, b
          c}, ch{} {}
   bool has_chd() const { return ch[0] !=
        nullptr; }
   bool contains(int q) const {
     F3 if (ori(v[p[i]], v[p[R(i)]], v[q])
           < 0)
        return false;
     return true;
   bool check(int q) const {
     return in_cc({v[p[0]], v[p[1]], v[p
[2]]}, v[q]); }
} pool[maxn * 10], *it, *root;
/* SPLIT_HASH_HERE */
 void link(const E &a, const E &b) {
   if (a.t) a.t->e[a.side] = b;
   if (b.t) b.t->e[b.side] = a;
void flip(Tri *A, int a) {
  auto [B, b] = A->e[a]; /* flip edge
    between A,B */
   if (!B || !A->check(B->p[b])) return;
   Tri *X = new (it++) Tri(A->p[R(a)], B->
        p[b], A \rightarrow p[a];
   Tri *Y = new (it++) Tri(B \rightarrow p[R(b)], A \rightarrow
        p[a], B->p[b]);
   link(E(X, 0), E(Y, 0));
   link(E(X, 1), A\rightarrow e[L(a)]); link(E(X, 2))
          B\rightarrow e[R(b)];
   link(E(Y, 1), B\rightarrow e[L(b)]); link(E(Y, 2))
         , A->e[R(a)]);
   A \rightarrow ch = B \rightarrow ch = \{X, Y, nullptr\};
   flip(X, 1); flip(X, 2); flip(Y, 1);
        flip(Y, 2);
 void add_point(int p) {
  Tri *r = root;
   while (r->has_chd()) for (Tri *c: r->ch
     if (c \&\& c\rightarrow contains(p)) \{ r = c;
          break; }
   array<Tri*, 3> t; /* split into 3
        triangles */
   F3 t[i] = new (it++) Tri(r->p[i], r->p[
  R(i)], p);
F3 link(E(t[i], 0), E(t[R(i)], 1));
   F3 link(E(t[i], 2), r->e[L(i)]);
   r->ch = t;
   F3 flip(t[i], 2);
auto build(const vector<P> &p) {
   it = pool; int n = (int)p.size();
   vector<int> ord(n); iota(all(ord), 0);
   shuffle(all(ord), mt19937(114514));
   root = new (it++) Tri(n, n + 1, n + 2);
   copy_n(p.data(), n, v); v[n++] = P(-C,
        -C);
  v[n++] = P(C * 2, -C); v[n++] = P(-C, C * 2);
   for (int i : ord) add_point(i);
   vector<array<int, 3>> res;
   for (Tri *now = pool; now != it; now++)
     if (!now->has_chd()) res.push_back(
          now->p);
   return res;
| }
```

8.5 DirInPoly

```
| bool DIP(const auto &p, int i, P dir) {
| const int n = (int)p.size();
| P A = p[i+1==n ? 0 : i+1] - p[i];
| P B = p[i==0 ? n-1 : i-1] - p[i];
| if (auto C = cross(A, B); C < 0)
| return cross(A, dir) >= 0 || cross(
| dir, B) >= 0;
| else
| return cross(A, dir) >= 0 && cross(
| dir, B) >= 0;
| // is Seg(p[i], p[i]+dir*eps) in p? (
| non-strict)
| // p is counterclockwise simple polygon
```

8.6 FarthestPair

```
|// p is CCW convex hull w/o colinear
    points
int n = (int)p.size(), pos = 1; lld ans =
    0;
for (int i = 0; i < n; i++) {
    P e = p[(i + 1) % n] - p[i];
    while (cross(e, p[(pos + 1) % n] - p[i
        ]) >
        cross(e, p[pos] - p[i]))
    pos = (pos + 1) % n;
    for (int j: {i, (i + 1) % n})
        ans = max(ans, norm(p[pos] - p[j]));
} // tested @ AOJ CGL_4_B
```

8.7 HPIGeneralLine

```
struct Line {
   lld a, b, \bar{c}; // ax + by + c <= 0
   P dir() const { return P(a, b); }
   Line(lld ta, lld tb, lld tc) : a(ta), b
        (tb), c(tc) {}
   Line(P S, P T):a(IM(T-S)),b(-RE(T-S)),c
        (cross(T,S)) {}
 }; using LN = const Line &;
PF intersect(LN A, LN B) {
  llf c = cross(A.dir(), B.dir());
i128 a = i128(A.c) * B.a - i128(B.c) *
       A.a;
   i128 b = i128(A.c) * B.b - i128(B.c) *
        A.b;
   return PF(-b / c, a / c);
 bool cov(LN 1, LN A, LN B) {
   i128 c = cross(A.dir(), B.dir());
i128 a = i128(A.c) * B.a - i128(B.c) *
        A.a;
   i128 b = i128(A.c) * B.b - i128(B.c) *
        A.b;
   return sgn(a * l.b - b * l.a + c * l.c)
         * sgn(c) >= 0;
bool operator<(LN a, LN b) {</pre>
   if (int c = argCmp(a.dir(), b.dir()))
    return c == -1;
   return i128(abs(b.a) + abs(b.b)) * a.c
                      i128(abs(a.a) + abs(a.b)
                           )) * b.c;
i }
```

8.8 HalfPlaneIntersection

```
struct Line {
  P st, ed, dir;
  Line (P s, P e) : st(s), ed(e), dir(e -
        s) {}
}; using LN = const Line &;
PF intersect(LN A, LN B) {
  llf t = cross(B.st - A.st, B.dir) /
    llf(cross(A.dir, B.dir));
  return toPF(A.st) + toPF(A.dir) * t; //
        C^3 / C^2
bool cov(LN 1, LN A, LN B) {
  i128 u = cross(B.st-A.st, B.dir);
  i128 v = cross(A.dir, B.dir);
  // ori(l.st, l.ed, A.st + A.dir*(u/v))
       <= 0?
  i128 \times = RE(A.dir) * u + RE(A.st - l.st
      ) * v;
  i128 y = IM(A.dir) * u + IM(A.st - l.st
       ) * v;
  return sgn(x*IM(l.dir) - y*RE(l.dir)) *
        sgn(v) >= 0;
} // x, y are C^3, also sgn<i128> is
    needed
bool operator<(LN a, LN b) {</pre>
  if (int c = argCmp(a.dir, b.dir))
    return c == -1;
  return ori(a.st, a.ed, b.st) < 0;</pre>
// cross(pt-line.st, line.dir)<=0 <-> pt
     in half plane
   the half plane is the LHS when going
     from st to ed
llf HPI(vector<Line> &q) {
  sort(q.begin(), q.end());
```

```
int n = (int)q.size(), l = 0, r = -1;
for (int i = 0; i < n; i++) {
     if (i && !argCmp(q[i].dir, q[i-1].dir
          )) continue;
     while (l < r && cov(q[i], q[r-1], q[r
          ])) --r;
     while (1 < r && cov(q[i], q[l], q[l
          +1])) ++1;
    q[++r] = q[i];
  while (l < r && cov(q[l], q[r-1], q[r])
  while (l < r && cov(q[r], q[l], q[l+1])</pre>
       ) ++1;
  n = r - l + 1; // q[l .. r] are the
       lines
  if (n \le 2 \mid | argCmp(q[l].dir, q[r].
        dir)) return 0;
  vector<PF> pt(n);
  for (int i = 0; i < n; i++)
    pt[i] = intersect(q[i+l], q[(i+1)%n+l])
          ]);
  return area(pt);
} // test @ 2020 Nordic NCPC : BigBrother
```

8.9 HullCut

8.10 KDTree

```
struct KDTree {
  struct Node {
    int x, y, x1, y1, x2, y2, id, f; Node
*L, *R;
  } tree[maxn], *root;
  illd dis2(int x1, int y1, int x2, int y2
    lld dx = x1 - x2, dy = y1 - y2;
return dx * dx + dy * dy;
  static bool cmpx(Node& a, Node& b) {
       return a.x<b.x; }</pre>
  static bool cmpy(Node& a, Node& b) {
       return a.y<b.y; }</pre>
  void init(vector<pair<int,int>> &ip) {
    for (int i = 0; i < ssize(ip); i++)</pre>
      tie(tree[i].x, tree[i].y) = ip[i],
            tree[i].id = i;
    root = build(0, (int)ip.size()-1, 0);
  Node* build(int L, int R, int d) {
    if (L>R) return nullptr;
    int M = (L+R)/2;
    nth_element(tree+L,tree+M,tree+R+1,d
         %2?cmpy:cmpx);
    Node &o = tree[M]; o.f = d \% 2;
    0.x1 = 0.x2 = \overline{0.x}; 0.y1 = 0.y2 = 0.y
    o.L = build(L, M-1, d+1); o.R = build
         (M+1, R, d+1);
    for (Node *s: {o.L, o.R}) if (s) {
    o.x1 = min(o.x1, s->x1); o.x2 = max
            (0.x2, s->x2);
      o.y1 = min(o.y1, s->y1); o.y2 = max
            (o.y2, s->y2);
    return tree+M;
  bool touch(int x, int y, lld d2, Node *
       r){
    lld d = (lld) sqrt(d2) + 1
```

return x >= r->x1 - d && x <= r->x2 + d &&

8.11 MinMaxEnclosingRect

```
// from 8BQube, plz ensure p is strict
    convex hull
  const llf INF = 1e18, qi = acos(-1) / 2 *
  pair<llf, llf> solve(const vector<P> &p)
          llf mx = 0, mn = INF; int n = (int)p.
                           size();
          for (int i = 0, u = 1, r = 1, l = 1; i
                              < n; ++i) {
  #define Z(v) (p[(v) % n] - p[i])
                 P e = Z(i + 1);
                 while (cross(e, Z(u + 1)) > cross(e,
                                    Z(u))) ++u;
                  while (dot(e, Z(r + 1)) > dot(e, Z(r))
                                   )) ++r;
                   if(!i) l = r + 1;
                 while (dot(e, Z(l + 1)) < dot(e, Z(l))
                                   )) ++l;
                 P D = p[r \% n] - p[l \% n];
                 llf H = cross(e, Z(u)) / llf(norm(e))
                 mn = min(mn, dot(e, D) * H);
                 llf B = sqrt(norm(D)) * sqrt(norm(Z(u))) * sqrt(n
                                    )));
                 mx = max(mx, B * sin(deg) * sin(deg))
          return {mn, mx};
} // test @ UVA 819
```

8.12 MinkowskiSum

```
// A, B are strict convex hull rotate to
     min by (X, Y)
vector<P> Minkowski(vector<P> A, vector<P
     > B) {
   const int N = (int)A.size(), M = (int)B
        .size();
  vector<P> sa(N), sb(M), C(N + M + 1);
for (int i = 0; i < N; i++) sa[i] = A[(
        i+1)%N]-A[i];
   for (int i = 0; i < M; i++) sb[i] = B[(
        i+1)%M]-B[i];
   C[0] = A[0] + B[0];
  for (int i = 0, j = 0; i < N \mid I \mid j < M;
        ) {
     P = (j>=M \mid i \in N \& cross(sa[i],
          sb[j])>=0))
       ? sa[i++] : sb[j++];
    C[i + j] = e;
  partial_sum(all(C), C.begin()); C.
       pop_back();
  return convex_hull(C); // just to
       remove colinear
} // be careful if min(|A|,|B|)<=2</pre>
```

8.13 PointInHull

```
|bool isAnti(P a, P b) {
   return cross(a, b) == 0 && dot(a, b) <=</pre>
        0; }
 bool PIH(const vector<P> &h, P z, bool
     strict = true) {
   int n = (int)h.size(), a = 1, b = n -
       1, r = !strict;
   if (n < 3) return r && isAnti(h[0] - z,
        h[n-1] - z);
   if (ori(h[0],h[a],h[b]) > 0) swap(a, b)
   if (ori(h[0],h[a],z) >= r || ori(h[0],h
        [b],z) <= -r)
     return false;
   while (abs(a - b) > 1) {
     int c = (a + b) / 2;
     (ori(h[0], h[c], z) > 0 ? b : a) = c;
   return ori(h[a], h[b], z) < r;</pre>
```

8.14 PointInPoly

```
| bool PIP(const vector<P> &p, P z, bool strict = true) {
| int cnt = 0, n = (int)p.size();
| for (int i = 0; i < n; i++) {
| P A = p[i], B = p[(i + 1) % n];
| if (isInter(Seg(A, B), z)) return !
| strict;
| auto zy = IM(z), Ay = IM(A), By = IM(B);
| cnt ^= ((zy<Ay) - (zy<By)) * ori(z, A, B) > 0;
| }
| return cnt;
```

8.15 PointInPolyFast

```
vector<int> PIPfast(vector<P> p, vector<P
    > q) {
  const int N = int(p.size()), Q = int(q.
      size());
  vector<pair<P, int>> evt; vector<Seg>
      edge;
  for (int i = 0; i < N; i++) {
    int a = i, b = (i + 1) \% N;
    P A = p[a], B = p[b];
    assert (A < B | \overline{B} < A); // std::
         operator<
    if (B < A) swap(A, B);
    evt.emplace_back(A, i); evt.
         emplace_back(B, ~i);
    edge.emplace_back(A, B);
  for (int i = 0; i < 0; i++)
    evt.emplace_back(q[i], i + N);
  sort(all(evt));
 auto vtx = p; sort(all(vtx));
auto eval = [](const Seg &a, lld x) ->
      11f {
    if (RE(a.dir) == 0) {
      assert (x == RE(a.st));
      return IM(a.st) + llf(IM(a.dir)) /
    llf t = (x - RE(a.st)) / llf(RE(a.dir))
        ));
    return IM(a.st) + IM(a.dir) * t;
 11d cur_x = 0;
  auto cmp = [&](const Seg &a, const Seg
      &b) -> bool {
    if (int s = sgn(eval(a, cur_x) - eval
         (b, cur_x)))
      return s == -1; // be careful: sgn<
          llf>, sgn<lld>
    int s = sgn(cross(b.dir, a.dir));
    if (cur_x != RE(a.st) && cur_x != RE(
         b.st)) s *= -1;
    return s == -1;
  namespace pbds = __gnu_pbds;
 pbds::tree<Seg, int, decltype(cmp),</pre>
    pbds::rb_tree_tag,
```

```
pbds::
          tree_order_statistics_node_update
          > st(cmp);
   auto answer = [\&](P ep) {
     if (binary_search(all(vtx), ep))
       return 1; // on vertex
     Seg H(ep, ep); // ??
     auto it = st.lower_bound(H);
     if (it != st.end() && isInter(it->
        first, ep))
return 1; // on edge
     if (it != st.begin() && isInter(prev())
          it)->first, ep))
       return 1; // on edge
     auto rk = st.order_of_key(H);
return rk % 2 == 0 ? 0 : 2; // 0:
          outside, 2: inside
   vector<int> ans(Q);
   for (auto [ep, i] : evt) {
     cur_x = RE(ep);
if (i < 0) { // remove</pre>
       st.erase(edge[~i]);
     } else if (i < N) { // insert</pre>
       auto [it, succ] = st.insert({edge[i
            ], i});
       assert(succ);
     } else ans[i - N] = answer(ep);
   return ans;
|} // test @ AOJ CGL_3_C
```

8.16 PolyUnion

```
llf polyUnion(const vector<vector<P>> &p)
  vector<tuple<P, P, int>> seg;
  for (int i = 0; i < ssize(p); i++)
    for (int j = 0, m = int(p[i].size());
         j < m; j++)
      seg.emplace_back(p[i][j], p[i][(j +
           1) % m], i);
  llf ret = 0; // area of p[i] must be
      non-negative
  for (auto [A, B, i] : seg) {
    vector<pair<llf, int>> evt{{0, 0},
        {1, 0}};
    for (auto [C, D, j] : seg) {
      int sc = ori(A, B, C), sd = ori(A,
          B, D);
      if (sc != sd && i != j && min(sc,
           sd) < 0) {
        llf sa = cross(D-C, A-C), sb =
            cross(D-C, B-C);
        evt.emplace_back(sa / (sa - sb),
            sgn(sc - sd));
      } else if (!sc && !sd && j < i</pre>
          && sgn(dot(B - A, D - C)) > 0)
        evt.emplace_back(real((C - A) / (
            B - A)), 1);
        evt.emplace_back(real((D - A) / (
            B - A)), -1);
     }
    sort(evt.begin(), evt.end());
    llf sum = 0, last = 0; int cnt = 0;
    for (auto [q, c] : evt) {
      if (!cnt) sum += q - last;
cnt += c; last = q;
   ret += cross(A, B) * sum;
  return ret / 2;
```

8.17 RotatingSweepLine

```
void rotatingSweepLine(const vector<P> &p
      ) {
   const int n = int(p.size());
   vector<Event> e; e.reserve(n * (n - 1)
        / 2);
  for (int i = 0; i < n; i++)
  for (int j = i + 1; j < n; j++)
    e.emplace_back(makePositive(p[i] -</pre>
             p[j]), i, j);
   sort(all(e));
   vector<int> ord(n), pos(n);
   iota(all(ord), 0);
   sort(all(ord), [&p](int i, int j) {
     return cmpxy(p[i], p[j]); });
   for (int i = 0; i < n; i++) pos[ord[i]]
         = i;
   const auto makeReverse = [](auto &v) {
     sort(all(v)); v.erase(unique(all(v)),
           v.end());
     vector<pair<int,int>> segs;
     for (size_t i = 0, j = 0; i < v.size
          (); i = j) {
       for (; j < v.size() && v[j] - v[i]</pre>
             <= j - i; j++);
       segs.emplace_back(v[i], v[j - 1] +
            1 + 1);
     return segs;
   for (size_t i = 0, j = 0; i < e.size();</pre>
     i = j) {
/* do here */
     vector<size_t> tmp;
     for (; j < e.size() && !(e[i] < e[j])</pre>
           ; j++)
       tmp.push_back(min(pos[e[j].u], pos[
             e[j].v]));
     for (auto [l, r] : makeReverse(tmp))
       reverse(ord.begin() + 1, ord.begin
            () + r);
       for (int t = 1; t < r; t++) pos[ord</pre>
             [t]] = t;
  }
|}
```

8.18 SegIsIntersect

```
struct Seg { // closed segment
 P st, dir; // represent st + t*dir for 0<=t<=1
  Seg(P s, P e) : st(s), dir(e - s) {}
 static bool valid(lld p, lld q) {
    // is there t s.t. 0 <= t <= 1 && qt
         == p ?
   if (q < 0) q = -q, p = -p;
return sgn(0 - p) \le 0 \& sgn(p - q)
 vector<P> ends() const { return { st,
       st + dir }; }
template <typename T> bool isInter(T A, P
     p) {
 if (sgn(norm(A.dir)) == 0)
    return sgn(norm(p - A.st)) == 0; //
BE CAREFUL
  return sgn(cross(p - A.st, A.dir)) == 0
        &&
    T::valid(dot(p - A.st, A.dir), norm(A
         .dir));
template <typename U, typename V>
bool isInter(U A, V B) {
 if (sgn(cross(A.dir, B.dir)) == 0) { //
BE (AREFUL
    bool res = false;
    for (P p: A.ends()) res |= isInter(B,
          p);
    for (P p: B.ends()) res |= isInter(A,
          p);
    return res;
 P D = B.st - A.st; lld C = cross(A.dir,
        B.dir):
  return U::valid(cross(D, B.dir), C) &&
```

```
V::valid(cross(D, A.dir), C);
```

8.19 SegSegDist

```
// be careful of abs<complex<int>> (
     replace _abs below)
llf PointSegDist(P A, Seg B) {
  if (B.dir == P(0)) return _abs(A - B.st
       );
  if (sgn(dot(A - B.st, B.dir)) *
    sgn(dot(A - B.ed, B.dir)) <= 0)
    return abs(cross(A - B.st, B.dir)) /</pre>
          _abs(B.dir);
  return min(_abs(A - B.st), _abs(A - B.
        ed));
llf SegSegDist(const Seg &s1, const Seg &
     52) {
  if (isInter(s1, s2)) return 0;
  return min({
       PointSegDist(s1.st, s2),
       PointSegDist(s1.ed, s2),
       PointSegDist(s2.st, s1);
       PointSegDist(s2.ed, s1) });
} // test @ QOJ2444 / PTZ19 Summer.D3
```

8.20 SimulateAnnealing

```
11f anneal() {
  mt19937 rnd_engine(seed);
  uniform_real_distribution<llf> rnd(0,
       1);
  const llf dT = 0.001;
   // Argument p
  llf S_cur = calc(p), S_best = S_cur;
for (llf T = 2000; T > EPS; T -= dT) {
    // Modify p to p_prime
    const llf S_prime = calc(p_prime);
     const llf delta_c = S_prime - S_cur;
    llf prob = min((llf)1, exp(-delta_c /
           T));
    if (rnd(rnd_engine) <= prob)</pre>
      S_cur = S_prime, p = p_prime;
     if (S_prime < S_best) // find min</pre>
      S_best = S_prime, p_best = p_prime;
  return S_best;
```

8.21 TangentPointToHull

8.22 TriCenter

```
| 0 = ... // see min circle cover
| G = (A + B + C) / 3;
| H = G * 3 - 0 * 2; // orthogonal center
| llf a = abs(B - C), b = abs(A - C), c =
| abs(A - B);
| I = (a * A + b * B + c * C) / (a + b + c)
| ;
| // FermatPoint: minimizes sum of distance
| // if max. angle >= 120 deg then vertex
| // otherwise, make eq. triangle AB'C, CA'
| B, BC'A
```

8.23 Voronoi

```
void build_voronoi_cells(auto &&p, auto
     &&res) {
  vector<vector<int>> adj(p.size());
   for (auto f: res) F3 {
     int a = f[i], b = f[R(i)];
     if (a \ge p.size() \mid | b \ge p.size())
         continue;
    adj[a].emplace_back(b);
  // use `adj` and `p` and HPI to build
       cells
   for (size_t i = 0; i < p.size(); i++) {</pre>
     vector<Line> ls = frame; // the frame
     for (int j : adj[i]) {
      P m = p[i] + p[j], d = rot90(p[j] -
            p[i]);
       assert (norm(d) != 0);
       ls.emplace_back(m, m + d); //
           doubled coordinate
    } // HPI(ls)
  }
}
```

9 Miscellaneous

9.1 Cactus 1

```
auto work = [&](const vector<int> cycle)
   // merge cycle info to u?
  int len = cycle.size(), u = cycle[0];
auto dfs = [&](auto dfs, int u, int p) {
  par[u] = p;
   vis[u] = 1;
  for (auto v : adj[u]) {
    if (v == p) { continue; }
    if (vis[v] == 0) {
      dfs(dfs, v, u);
       if (!cyc[v]) { // merge dp }
    } else if (vis[v] == 1) {
      for (int w = u; w != v; w = par[w])
        cyc[w] = 1;
    } else {
      vector<int> cycle = {u};
      for (int w = v; w != u; w = par[w])
            { cycle.push_back(w); }
      work(cycle);
  vis[u] = 2;
|};
```

9.2 Cactus 2

```
// a component contains no articulation
     point, so P2 is a component
// but not a vertex biconnected component
      by definition
// resulting bct is rooted
struct BlockCutTree {
  int n, square = 0, cur = 0;
  vector<int> low, dfn, stk;
  vector<vector<int>> adj, bct;
  BlockCutTree(int n) : n(n), low(n), dfn
  (n, -1), adj(n), bct(n) {}
void build() { dfs(0); }
void addEdge(int u, int v) { adj[u].
       push_back(v), adj[v].push_back(u);
  void dfs(int u) {
    low[u] = dfn[u] = cur++;
    stk.push_back(u);
    for (auto v : adj[u]) {
      if (dfn[v] == -1) {
        dfs(v);
         low[u] = min(low[u], low[v]);
        if (low[v] == dfn[u]) {
           bct.emplace_back();
           int x;
```

```
do {
             x = stk.back();
             stk.pop_back();
             bct.back().push_back(x);
           } while (x != v);
           bct[u].push_back(n + square);
           square++;
        }
      } else {
        low[u] = min(low[u], dfn[v]);
  }
|};
```

Dancing Links

```
#include <bits/stdc++.h>
using namespace std;
// tioj 1333
])
const int NN = 40000, RR = 200;
template<bool E> // E: Exact, NN: num of
    1s, RR: num of rows
struct DLX {
  int lt[NN], rg[NN], up[NN], dn[NN], rw[
       NN], cl[NN], bt[NN], s[NN], head,
       sz, ans;
  int rows, columns;
  bool vis[NN];
  bitset<RR> sol, cur; // not sure
  void remove(int c) {
    if (E) lt[rg[c]] = lt[c], rg[lt[c]] =
          rg[c];
    TRAV(i, dn, c) {
   if (E) {
        TRAV(j, rg, i)
up[dn[j]] = up[j], dn[up[j]] =
               dn[j], --s[cl[j]];
      } else {
        lt[rg[i]] = lt[i], rg[lt[i]] = rg
             [i];
    }
  void restore(int c) {
    TRAV(i, up, c) {
      if (E) {
        TRAV(j, lt, i)
          ++s[cl[j]], up[dn[j]] = j, dn[
               up[j]] = j;
      } else {
        lt[rg[i]] = rg[lt[i]] = i;
    if (E) lt[rg[c]] = c, rg[lt[c]] = c;
  void init(int c) {
    rows = 0, columns = c;
    for (int i = 0; i < c; ++i) {
    up[i] = dn[i] = bt[i] = i;
      lt[i] = i == 0 ? c : i - 1;
rg[i] = i == c - 1 ? c : i + 1;
      s[i] = 0;
    rg[c] = 0, lt[c] = c - 1;
    up[c] = dn[c] = -1;
head = c, sz = c + 1;
  void insert(const vector<int> &col) {
    if (col.empty()) return;
    int f = sz;
    for (int i = 0; i < (int)col.size();</pre>
         ++i) {
      int c = col[i], v = sz++;
      dn[bt[c]] = v;
      up[v] = bt[c], bt[c] = v;
      rg[v] = (i + 1 == (int)col.size() ?
            f : v + 1:
      rw[v] = rows, cl[v] = c;
      ++s[c];
      if (i > 0) lt[v] = v - 1;
    ++rows, lt[f] = sz - 1;
```

```
int h() {
    int ret = 0;
    fill_n(vis, sz, false);
TRAV(x, rg, head) {
      if (vis[x]) continue;
      vis[x] = true, ++ret;
TRAV(i, dn, x) TRAV(j, rg, i) vis[
           cl[j]] = true;
    return ret;
 void dfs(int dep) {
    if (dep + (E ? 0 : h()) >= ans)
         return:
    if (rg[head] == head) return sol =
         cur, ans = dep, void();
    if (dn[rg[head]] == rg[head]) return;
    int w = rg[head];
    TRAV(x, rg, head) if (s[x] < s[w]) w
         = x;
    if (E) remove(w);
    TRAV(i, dn, w) {
      if (!E) remove(i);
      TRAV(j, rg, i) remove(E ? cl[j] : j
      cur.set(rw[i]), dfs(dep + 1), cur.
           reset(rw[i]);
      TRAV(j, lt, i) restore(E ? cl[j] :
           j);
      if (!E) restore(i);
    if (E) restore(w);
  int solve() {
    for (int i = 0; i < columns; ++i)</pre>
      dn[bt[i]] = i, up[i] = bt[i];
    ans = 1e9, sol.reset(), dfs(0);
    return ans:
 }
int main() {
    int n, m; cin >> n >> m;
DLX<true> solver;
    solver.init(m);
    for (int i = 0; i < n; i++){
        vector<int> add;
        for (int j = 0; j < m; j++){
             int x; cin >> x;
            if (x == 1) {
                 add.push_back(j);
        solver.insert(add);
    cout << solver.solve() << '\n';</pre>
    return 0;
9.4 Offline Dynamic MST
```

```
int cnt[maxn], cost[maxn], st[maxn], ed[
    maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed
     , qr[i].second = weight after
     operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v
     contains edges i such that cnt[i] ==
void contract(int 1, int r, vector<int> v
      vector<int> &x, vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int
        } (i
      if (cost[i] == cost[j]) return i <</pre>
           j;
      return cost[i] < cost[j];</pre>
      }):
  djs.save();
  for (int i = l; i <= r; ++i) djs.merge(</pre>
       st[qr[i].first], ed[qr[i].first]);
  for (int i = 0; i < (int)v.size(); ++i)</pre>
       (djs.find(st[v[i]]) != djs.find(ed
         [v[i]])) {
      x.push_back(v[i]);
```

```
\label{eq:discontinuity} \mbox{djs.merge(st[v[i]], ed[v[i]]);}
   }
 djs.undo();
 djs.save();
 for (int i = 0; i < (int)x.size(); ++i)</pre>
        djs.merge(st[x[i]], ed[x[i]]);
  for (int i = 0; i < (int)v.size(); ++i)</pre>
    if (djs.find(st[v[i]]) != djs.find(ed
         [v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
 dis.undo();
void solve(int l, int r, vector<int> v,
     long long c) {
  if (l == r) {
    cost[qr[l].first] = qr[l].second;
    if (st[qr[l].first] == ed[qr[l].first]
         1) {
      printf("%lld\n", c);
      return;
    int minv = qr[l].second;
    for (int i = 0; i < (int)v.size(); ++
         i) minv = min(minv, cost[v[i]]);
    printf("%lld\n", c + minv);
    return;
 int m = (l + r) >> 1;
 vector<int> lv = v, rv = v;
 vector<int> x, y;
  for (int i = m + 1; i \le r; ++i) {
    cnt[qr[i].first]--
    if (cnt[qr[i].first] == 0) lv.
         push_back(qr[i].first);
 contract(l, m, lv, x, y);
long long lc = c, rc = c;
 djs.save();
  for (int i = 0; i < (int)x.size(); ++i)</pre>
    lc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
 solve(l, m, y, lc);
 djs.undo();
 x.clear(), y.clear();
for (int i = m + 1; i <= r; ++i) cnt[qr
       [i].first]++;
  for (int i = l; i <= m; ++i) {
    cnt[qr[i].first]--
    if (cnt[qr[i].first] == 0) rv.
         push_back(qr[i].first);
  contract(m + 1, r, rv, x, y);
 djs.save();
  for (int i = 0; i < (int)x.size(); ++i)
    rc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
 solve(m + 1, r, y, rc);
 djs.undo();
 for (int i = l; i <= m; ++i) cnt[qr[i].</pre>
       first]++;
```

Matroid Intersection 9.5

```
• x \to y if S - \{x\} \cup \{y\} \in I_1 with cost(\{y\}).
• source \rightarrow y \text{ if } S \cup \{y\} \in I_1 \text{ with } cost(\{y\}).
• y \to x if S - \{x\} \cup \{y\} \in I_2 with -cost(\{y\}).
• y \to sink if S \cup \{y\} \in I_2 with -cost(\{y\}).
```

Augmenting path is shortest path from source to

9.6Euler Tour

```
vector<int> euler, vis(V);
auto dfs = [&](auto dfs, int u) -> void {
 while (!adj[u].empty()) {
```

```
struct SegmentTree {
  int n;
  struct node {
    i64 mx1, mx2, mxc;
    i64 mn1, mn2, mnc;
    i64 add:
    i64 sum;
    node(i64 v = 0) {
       mx1 = mn1 = sum = v;
       mxc = mnc = 1;
       add = 0;
      mx2 = -9e18, mn2 = 9e18;
    }
  vector<node> t;
  // build
  void push(int id, int l, int r) {
    auto& c = t[id];
    int m = l + r \gg 1;
    if (c.add != 0) {
       apply_add(id << 1, 1, m, c.add);</pre>
       apply_add(id \ll 1 | 1, m + 1, r, c.
           add):
      c.add = 0;
    apply_min(id << 1, 1, m, c.mn1);</pre>
    apply_min(id \ll 1 | 1, m + 1, r, c.
    apply_max(id << 1, 1, m, c.mx1);
apply_max(id << 1 | 1, m + 1, r, c.</pre>
         mx1);
  void apply_add(int id, int l, int r,
       i64 v) {
     if (v == 0) {
      return;
    auto& c = t[id];
    c.add += v;
c.sum += v * (r - l + 1);
    c.mx1 += v;
    c.mn1 += v;
    if (c.mx2 != -9e18) {
       c.mx2 += v;
    if (c.mn2 != 9e18) {
       c.mn2 += v;
  void apply_min(int id, int l, int r,
       i64 v) {
    auto& c = t[id];
    if (v \leftarrow c.mn1) {
       return;
    c.sum -= c.mn1 * c.mnc;
    c.mn1 = v;
c.sum += c.mn1 * c.mnc;
    if (l == r \mid \mid v >= c.mx1) {
      c.mx1 = v;
    } else if (v > c.mx2) {
       c.mx2 = v;
    }
  void apply_max(int id, int l, int r,
       i64 v) {
    auto& c = t[id];
     if (v \ge c.mx1) {
       return:
    c.sum -= c.mx1 * c.mxc;
```

```
c.mx1 = v:
  c.sum += c.mx1 * c.mxc;
  if (1 == r || v <= c.mn1) {
    c.mn1 = v;
  } else if (v < c.mn2) {
   c.mn2 = v;</pre>
  }
void pull(int id) {
  auto &c = t[id], &lc = t[id << 1], &
       rc = t[id << 1 | 1];
  c.sum = lc.sum + rc.sum;
  if (lc.mn1 == rc.mn1) {
    c.mn1 = lc.mn1;
    c.mn2 = min(lc.mn2, rc.mn2);
c.mnc = lc.mnc + rc.mnc;
  } else if (lc.mn1 < rc.mn1) {
   c.mn1 = lc.mn1;</pre>
    c.mn2 = min(lc.mn2, rc.mn1);
    c.mnc = lc.mnc;
  } else {
    c.mn1 = rc.mn1;
c.mn2 = min(lc.mn1, rc.mn2);
    c.mnc = rc.mnc;
  if (lc.mx1 == rc.mx1) {
    c.mx1 = lc.mx1;
    c.mx2 = max(lc.mx2, rc.mx2);
c.mxc = lc.mxc + rc.mxc;
  } else if (lc.mx1 > rc.mx1) {
   c.mx1 = lc.mx1;
    c.mx2 = max(lc.mx2, rc.mx1);
    c.mxc = lc.mxc;
  } else {
    c.mx1 = rc.mx1;
    c.mx2 = max(lc.mx1, rc.mx2);
    c.mxc = rc.mxc;
void range_chmin(int id, int 1, int r,
     int ql, int qr, i64 v) {
  if (r < ql \mid | l > qr \mid | v >= t[id].
       mx1) {
    return;
  if (ql \le l \& r \le qr \& v > t[id].
       mx2) {
    apply_max(id, l, r, v);
    return;
  push(id, l, r);
int m = l + r >> 1;
  range_chmin(id << 1, l, m, ql, qr, v)
  range_chmin(id \ll 1 | 1, m + 1, r, ql
         qr, v);
  pull(id);
void range_chmin(int ql, int qr, i64 v)
  range_chmin(1, 0, n - 1, ql, qr, v);
void range_chmax(int id, int l, int r,
  int ql, int qr, i64 v) {
if (r < ql || l > qr || v <= t[id].
       mn1) {
    return:
  if (ql \ll l \& r \ll qr \& v \ll t[id].
       mn2) {
    apply_min(id, l, r, v);
    return;
  push(id, l, r);
int m = l + r >> 1;
  range\_chmax(id << 1, l, m, ql, qr, v)
  range\_chmax(id << 1 \mid 1, m + 1, r, ql)
         qr, v);
  pull(id);
void range_chmax(int ql, int qr, i64 v)
  range_chmax(1, 0, n - 1, ql, qr, v);
void range_add(int id, int l, int r,
     int ql, int qr, i64 v) {
  if (r < ql || l > qr) {
```

```
return:
     if (ql <= l && r <= qr) {
        apply_add(id, l, r, v);
        return;
     push(id, l, r);
int m = l + r >> 1;
range_add(id << 1, l, m, ql, qr, v);
range_add(id << 1 | 1, m + 1, r, ql,</pre>
          qr, v);
     pull(id);
   }
   void range_add(int ql, int qr, i64 v) {
     range_add(1, 0, n - 1, ql, qr, v);
   i64 range_sum(int id, int l, int r, int
     ql, int qr) {
if (r < ql || l > qr) {
       return 0;
     if (ql \ll l \& r \ll qr) {
       return t[id].sum;
     push(id, 1, r);
     int m = l + r >> 1;
     return range_sum(id << 1, 1, m, q1,</pre>
           qr) + range_sum(id \ll 1 | 1, m +
            1, r, ql, qr);
   i64 range_sum(int ql, int qr) {
     return range_sum(1, 0, n - 1, ql, qr)
  }
};
```

9.8 Decimal

```
| from decimal import *
| setcontext(Context(prec=MAX_PREC, Emax=
| MAX_EMAX, rounding=ROUND_FLOOR))
| print(Decimal(input()) * Decimal(input())
| )
| from fractions import Fraction
| Fraction('3.14159').limit_denominator(10)
| .numerator # 22
```

9.9 AdaSimpson

```
template<typename Func, typename d =</pre>
     double>
struct Simpson {
  using pdd = pair<d, d>;
  pdd mix(pdd l, pdd r, optional<d> fm =
       {}) {
    dh = (r.X - 1.X) / 2, v = fm.
         value_or(f(1.X + h));
    return \{v, h / 3 * (1.Y + 4 * v + r.Y)\}
         )};
  d eval(pdd l, pdd r, d fm, d eps) {
  pdd m((1.X + r.X) / 2, fm);
  d s = mix(l, r, fm).second;
    auto [flm, sl] = mix(l, m);
    return eval(l, m, flm, eps / 2) +
       eval(m, r, fmr, eps / 2);
 d eval(d l, d r, d eps) {
    return eval({l, f(l)}, {r, f(r)}, f((
         1 + r) / 2, eps);
  d eval2(d l, d r, d eps, int k = 997) {
    d h = (r - l) / k, s = 0;
    con (data)
    for (int i = 0; i < k; ++i, l += h)
      s += eval(l, l + h, eps / k);
    return s;
 }
template<typename Func>
Simpson<Func> make_simpson(Func f) {
     return {f}; }
```

9.10 SB Tree

```
struct Q {
   ll p, q;
   Q go(Q b, ll d) { return {p + b.p*d, q
        + b.q*d}; }
 bool pred(Q);
 // returns smallest p/q in [lo, hi] such
 // pred(p/q) is true, and 0 \le p,q \le N
 Q frac_bs(ll N) {
   Q lo{0, 1}, hi{1, 0};
   if (pred(lo)) return lo;
   assert(pred(hi));
   bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
     for (int t = 0; t < 2 \&\& (t ? step/=2)
            : step*=2);)
       if (Q mid = hi.go(lo, len + step);
           mid.p > N \mid \mid mid.q > N \mid \mid dir \land
                  pred(mid))
       else len += step;
     swap(lo, hi = hi.go(lo, len));
     (dir ? L : H) = !!len;
   return dir ? hi : lo;
į }
```

9.11 Bitset LCS

```
| cin >> n >> m;
| for (int i = 1, x; i <= n; ++i)
| cin >> x, p[x].set(i);
| for (int i = 1, x; i <= m; i++) {
| cin >> x, (g = f) |= p[x];
| f.shiftLeftByOne(), f.set(0);
| ((f = g - f) ^= g) &= g;
| }
| cout << f.count() << '\n';</pre>
```

9.12 Hilbert Curve

9.13 Mo on Tree

```
void MoAlgoOnTree() {
 Dfs(0, -1);
  vector<int> euler(tk);
 for (int i = 0; i < n; ++i) {
  euler[tin[i]] = i;</pre>
    euler[tout[i]] = i;
 vector<int> l(q), r(q), qr(q), sp(q,
       -1);
  for (int i = 0; i < q; ++i) {
    if (tin[u[i]] > tin[v[i]]) swap(u[i],
         v[i]);
    int z = GetLCA(u[i], v[i]);
    sp[i] = z[i];
    if (z == u) l[i] = tin[u[i]], r[i] =
         tin[v[i]];
    else l[i] = tout[u[i]], r[i] = tin[v[
         i]];
    qr[i] = i;
 sort(qr.begin(), qr.end(), [&](int i,
       int j) {
      if (l[i] / kB == l[j] / kB) return
           r[i] < r[j];
      return l[i] / kB < l[j] / kB;</pre>
      });
 vector<bool> used(n);
```

9.14 N Queens

```
void solve(vector<int> &ret, int n) { //
     no sol when n=2,3
   if (n % 6 == 2) {
     for (int i = 2; i \leftarrow n; i += 2) ret.
          pb(i);
     ret.pb(3); ret.pb(1);
     for (int i = 7; i \leftarrow n; i += 2) ret.
          pb(i);
     ret.pb(5);
  } else if (n % 6 == 3) {
     for (int i = 4; i \le n; i += 2) ret.
          pb(i);
     ret.pb(2);
     for (int i = 5; i \le n; i += 2) ret.
          pb(i);
     ret.pb(1); ret.pb(3);
  } else {
     for (int i = 2; i <= n; i += 2) ret.
          pb(i);
     for (int i = 1; i <= n; i += 2) ret.
          pb(i);
  }
}
```

9.15 Rollback Mo

```
for (int l = 0, r = -1; auto [ql, qr, i]
      qs) ·
  if (ql / B == qr / B) {
    for (int j = ql; j <= qr; j++) {</pre>
       cntSmall[a[j]]++;
       ans[i] = max(ans[i], 1LL * b[a[j]]
     * cntSmall[a[j]]);
    for (int j = ql; j <= qr; j++) {
  cntSmall[a[j]]--;</pre>
    }
    continue;
  if (int block = ql / B; block != lst) {
    int x = min((block + 1) * B, n);
    while (r + 1 < x) \{ add(++r); \}
    while (r >= x) \{ del(r--); \}
    while (l < x) \{ del(l++); \}
    mx = 0:
    lst = block;
  while (r < qr) { add(++r); }</pre>
  i64 \text{ tmpMx} = mx;
  int tmpL = 1;
  while (l > ql) { add(--l); }
  ans[i] = mx;
  mx = tmpMx;
  while (l < tmpL) { del(l++); }</pre>
```

9.16 Subset Sum

```
template<size_t S> // sum(a) < S
bitset<S> SubsetSum(const int *a, int n)
   {
    vector<int> c(S);
    bitset<S> dp; dp[0] = 1;
    for (int i = 0; i < n; ++i) ++c[a[i ]];
    for (size_t i = 1; i < S; ++i) {
        while (c[i] > 2) c[i] -= 2, ++c[i * 2];
        while (c[i]--) dp |= dp << i;</pre>
```