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## Contents

## 1 Basic

#### 1.1 vimrc

```
se is nu ru et tgc sc hls cin cino+=j1 sw
     =4 sts=4 bs=2 mouse=a "encoding=utf
     -8 ls=2
syn on | colo desert | filetype indent on map <leader>b <ESC>:w<CR>:!g++ "%" -o "%<
     " -g -std=gnu++20 -DCKISEKI -Wall -
     Wextra -Wshadow -Wfatal-errors -
     Wconversion -fsanitize=address.
     undefined, float-divide-by-zero, float
success<CR>
map <leader>i <ESC>:!./"%<"<CR>
map <leader>r <ESC>:!cat 01.in && echo "
      ---" && ./"%<" < 01.in<CR>
map <leader>l :%d<bar>0r ~/t.cpp<CR>
ca Hash w !cpp -dD -P -fpreprocessed \|
    tr -d "[:space:]" \| md5sum \| cut -
     c-6
let c_no_curly_error=1
```

## 1.2 Default code

```
#include <bits/stdc++.h>
using namespace std;
using i64 = long long;
using ll = long long;
#define SZ(v) (ll)((v).size())
#define pb emplace_back
#define AI(i) begin(i), end(i)
#define X first
#define Y second
template<class T> bool chmin(T &a, T b) {
      return b < a && (a = b, true); }</pre>
template<class T> bool chmax(T &a, T b) {
      return a < b && (a = b, true); }</pre>
#ifdef KEV
#define DE(args...) kout("[ " + string(#
    args) + " ] = ", args)
void kout() { cerr << endl; }</pre>
template<class T, class ...U> void kout(T
      a, U ...b) { cerr << a << ' ', kout
     (b...); }
template<class T> void debug(T l, T r) {
     while (l != r) cerr << *l << " \n"[</pre>
     next(l)==r], ++l; }
#else
#define DE(...) 0
#define debug(...) 0
#endif
int main() {
  cin.tie(nullptr)->sync_with_stdio(false
  );
return 0;
```

#### 1.3 Fast Integer Input

```
|}
int readInt() {
   int x = 0;
   char c = get();
   while (!isdigit(c))
      c = get();
   while (isdigit(c)) {
      x = 10 * x + c - '0';
      c = get();
   }
   return x;
|}
```

### 1.4 Fast Python Input

```
|import sys, os, io
|input = io.BytesIO(os.read(0, os.fstat(0)
| .st_size)).readline
```

### 1.5 Pragma optimization

```
#pragma GCC optimize("Ofast", "no-stack-
    protector", "no-math-errno", "unroll
    -loops")
#pragma GCC target("sse,sse2,sse3,ssse3,
    sse4,sse4.2,popcnt,abm,mmx,avx,tune=
    native,arch=core-avx2,tune=core-avx2
    ")
#pragma GCC ivdep
```

## 2 Flows, Matching

#### 2.1 Flow

template <typename F>

```
struct Flow {
 static constexpr F INF = numeric_limits
      <F>::max() / 2;
 struct Edge {
   int to;
   F cap;
   Edge(int to, F cap) : to(to), cap(cap
        ) {}
 int n;
 vector<Edge> e;
 vector<vector<int>> adj;
 vector<int> cur, h;
 Flow(int n) : n(n), adj(n) {}
 bool bfs(int s, int t) {
   h.assign(n, -1);
   queue<int> q;
   h[s] = 0;
   q.push(s);
   while (!q.empty()) {
     int u = q.front();
     q.pop();
      for (int i : adj[u]) {
        auto [v, c] = e[i];
        if (c > 0 \& h[v] == -1) {
         h[v] = h[u] + 1;
         if (v == t) { return true; }
         q.push(v);
       }
     }
   return false;
 F dfs(int u, int t, F f) {
   if (u == t) { return f; }
   int j = adj[u][i];
     auto [v, c] = e[j];
if (c > 0 && h[v] == h[u] + 1) {
       Fa = dfs(v, t, min(r, c));
       e[i].cap -= a;
       e[j ^ 1].cap += a;
         -= a;
       if (r == 0) \{ return f; \}
     }
   }
   return f - r;
```

```
// can be bidirectional
void addEdge(int u, int v, F cf = INF,
     F cb = 0) {
  adj[u].push_back(e.size()), e.
       emplace_back(v, cf);
  adj[v].push_back(e.size()), e.
       emplace_back(u, cb);
F maxFlow(int s, int t) {
  F ans = 0;
  while (bfs(s, t)) {
    cur.assign(n, 0);
    ans += dfs(s, t, INF);
  return ans:
// do max flow first
vector<int> minCut() {
  vector<int> res(n);
  for (int i = 0; i < n; i++) { res[i]
= h[i] != -1; }
  return res;
}
```

#### 2.2 MCMF

```
template <class Flow, class Cost>
struct MinCostMaxFlow {
public:
  static constexpr Flow flowINF =
       numeric_limits<Flow>::max();
  static constexpr Cost costINF =
       numeric_limits<Cost>::max();
  MinCostMaxFlow() {}
  MinCostMaxFlow(int n) : n(n), g(n) {}
  int addEdge(int u, int v, Flow cap,
       Cost cost) {
    int m = int(pos.size());
    pos.push_back({u, int(g[u].size())});
    g[u].push_back({v, int(g[v].size()),
         cap, cost});
    g[v].push_back({u, int(g[u].size()) -
          1, 0, -cost});
    return m;
  struct edae {
    int u, v;
    Flow cap, flow;
Cost cost;
  }:
  edge getEdge(int i) {
    auto _e = g[pos[i].first][pos[i].
         second];
    auto _re = g[_e.v][_e.rev];
    return {pos[i].first, _e.v, _e.cap +
         _re.cap, _re.cap, _e.cost};
  vector<edge> edges() {
    int m = int(pos.size());
    vector<edge> result(m);
for (int i = 0; i < m; i++) { result[</pre>
         i] = getEdge(i); }
    return result;
  pair<Flow, Cost> maxFlow(int s, int t,
       Flow flow_limit = flowINF) {
       return slope(s, t, flow_limit).
       back(); }
  vector<pair<Flow, Cost>> slope(int s,
       int t, Flow flow_limit = flowINF)
    vector<Cost> dual(n, 0), dis(n);
    vector<int> pv(n), pe(n), vis(n);
auto dualRef = [&]() {
      fill(dis.begin(), dis.end(),
           costINF):
      fill(pv.begin(), pv.end(), -1);
      fill(pe.begin(), pe.end(), -1);
      fill(vis.begin(), vis.end(), false)
      struct Q {
        Cost key;
        int u:
        bool operator<(Q o) const {</pre>
             return key > o.key; }
```

```
priority_queue<Q> h;
       dis[s] = 0;
h.push({0, s});
       while (!h.empty()) {
          int u = h.top().u;
         h.pop();
         if (vis[u]) { continue; }
         vis[u] = true;
          if (u == t) { break; }
          for (int i = 0; i < int(g[u].size</pre>
              ()); i++) {
            auto e = g[u][i];
            if (vis[e.v] | l e.cap == 0)
                 continue;
            Cost cost = e.cost - dual[e.v]
                 + dual[u];
            if (dis[e.v] - dis[u] > cost) {
              dis[e.v] = dis[u] + cost;
              pv[e.v] = u;
              pe[e.v] = i;
              h.push({dis[e.v], e.v});
         }
       if (!vis[t]) { return false; }
       for (int v = 0; v < n; v++) {
         if (!vis[v]) continue;
         dual[v] -= dis[t] - dis[v];
       return true;
     Flow flow = 0;
     Cost cost = 0, prevCost = -1;
     vector<pair<Flow, Cost>> result;
     result.push_back({flow, cost});
while (flow < flow_limit) {</pre>
       if (!dualRef()) break;
       Flow c = flow_limit - flow;
       for (int v = t; v != s; v = pv[v])
         c = min(c, g[pv[v]][pe[v]].cap);
       for (int v = t; v != s; v = pv[v])
         auto& e = g[pv[v]][pe[v]];
         e.cap -= c:
         g[v][e.rev].cap += c;
       Cost d = -dual[s];
       flow += c;
cost += c * d;
       if (prevCost == d) { result.
            pop_back(); }
       result.push_back({flow, cost});
       prevCost = cost;
     return result;
  }
 private:
  int n:
   struct _edge {
     int v, rev;
Flow cap;
     Cost cost;
   vector<pair<int, int>> pos;
   vector<vector<_edge>> g;
1};
```

#### 2.3 GomoryHu Tree

```
auto gomory(int n, vector<array<int, 3>>
    e) {
  Flow<int, int> mf(n);
  for (auto [u, v, c] : e) { mf.addEdge(u
       , v, c, c); }
  vector<array<int, 3>> res;
  vector<int> p(n);
  for (int i = 1; i < n; i++) {
    for (int j = 0; j < int(e.size()); j</pre>
        ++) { mf.e[j << 1].cap = mf.e[j
         << 1 | 1].cap = e[j][2]; }
    int f = mf.maxFlow(i, p[i]);
    auto cut = mf.minCut();
    for (int j = i + 1; j < n; j++) { if
         (cut[i] == cut[j] && p[i] == p[j
         ]) { p[j] = i; }}
```

```
res.push_back({f, i, p[i]});
}
return res;
}
```

### 2.4 Global Minimum Cut

```
// 0(V ^ 3)
template <typename F>
struct GlobalMinCut {
  static constexpr int INF =
       numeric_limits<F>::max() / 2;
  int n:
  vector<int> vis, wei;
   vector<vector<int>> adj;
  GlobalMinCut(int n) : n(n), vis(n), wei
        (n), adj(n, vector<int>(n)) {}
   void addEdge(int u, int v, int w){
     adj[u][v] += w;
     adj[v][u] += w;
  int solve() {
     int sz = n;
     int res = INF, x = -1, y = -1;
     auto search = [&]() {
       fill(vis.begin(), vis.begin() + sz,
             0):
       fill(wei.begin(), wei.begin() + sz,
      x = y = -1;
       int mx, cur;
       for (int i = 0; i < sz; i++) {
         mx = -1, cur = 0;
         for (int j = 0; j < sz; j++) {
           if (wei[j] > mx) {
             mx = wei[j], cur = j;
         vis[cur] = 1, wei[cur] = -1;
         x = y;
y = cur;
         for (int j = 0; j < sz; j++) {
           if (!vis[j]) {
             wei[j] += adj[cur][j];
         }
       return mx;
     while (sz > 1) {
       res = min(res, search());
       for (int i = 0; i < sz; i++) {
         adj[x][i] += adj[y][i];
         adj[i][x] = adj[x][i];
       for (int i = 0; i < sz; i++) {
  adj[y][i] = adj[sz - 1][i];</pre>
         adj[i][y] = adj[i][sz - 1];
       sz--;
     return res;
  }
};
```

## 2.5 Bipartite Matching

```
struct BipartiteMatching {
  int n, m;
  vector<vector<int>> adj;
 vector<int> l, r, dis, cur;
BipartiteMatching(int n, int m) : n(n),
        m(m), adj(n), l(n, -1), r(m, -1),
        dis(n), cur(n) {}
  void addEdge(int u, int v) { adj[u].
       push_back(v); }
  void bfs() {
    vector<int> q;
    for (int u = 0; u < n; u++) {
      if (l[u] == -1) {
        q.push_back(u), dis[u] = 0;
      } else {
        dis[u] = -1;
    for (int i = 0; i < int(q.size()); i</pre>
         ++) {
```

```
int u = q[i];
for (auto_v : adj[u]) {
         if (r[v] != -1 && dis[r[v]] ==
              -1) {
           dis[r[v]] = dis[u] + 1;
           q.push_back(r[v]);
      }
  bool dfs(int u) {
     for (int \&i = cur[u]; i < int(adj[u].
          size()); i++) {
       int v = adj[u][i];
       if(r[v] == -1 \mid l \mid dis[r[v]] == dis[
            u] + 1 && dfs(r[v])) {
         l[u] = v, r[v] = u;
return true;
    }
     return false;
   int maxMatching() {
     int match = 0:
     while (true) {
       bfs();
       fill(cur.begin(), cur.end(), 0);
       int cnt = 0;
       for (int u = 0; u < n; u++) {
         if (l[u] == -1) {
           cnt += dfs(u);
       if (cnt == 0) {
         break;
       match += cnt;
     return match:
  auto minVertexCover() {
     vector<int> L, R;
     for (int u = 0; u < n; u++) {
       if (dis[u] == -1) {
         L.push_back(u);
       } else if (l[u] != -1) {
         R.push_back(l[u]);
     return pair(L, R);
};
```

#### 2.6 GeneralMatching

```
| struct GeneralMatching {
  int n:
  vector<vector<int>> adj;
  vector<int> match;
  GeneralMatching(int n) : n(n), adj(n),
       match(n, -1) {}
  void addEdge(int u, int v) {
    adj[u].push_back(v);
    adj[v].push_back(u);
  int maxMatching() {
    vector<int> vis(n), link(n), f(n),
         dep(n);
    auto find = [&](int u) {
      while (f[u] != u) \{ u = f[u] = f[f[
           u]]; }
    auto lca = [&](int u, int v) {
      u = find(u);
      v = find(v);
      while (u != v) {
         if (dep[u] < dep[v]) { swap(u, v)
        u = find(link[match[u]]);
      return u;
    queue<int> q;
    auto blossom = [&](int u, int v, int
      while (find(u) != p) {
```

```
link[u] = v
          v = match[u];
          if (vis[v] == 0) {
            vis[v] = 1;
            q.push(v);
          f[u] = f[v] = p;
          u = link[v];
     auto augment = [&](int u) {
       while (!q.empty()) { q.pop(); }
        iota(f.begin(), f.end(), 0);
fill(vis.begin(), vis.end(), -1);
q.push(u), vis[u] = 1, dep[u] = 0;
        while (!q.empty()){
          int u = q.front();
          q.pop();
          for (auto v : adj[u]) {
            if (vis[v] == -1) {
               vis[v] = 0;
              link[v] = u;
dep[v] = dep[u] + 1;
               if (match[v] == -1) {
                 for (int x = v, y = u, tmp;
y!= -1; x = tmp, y =
                        x == -1 ? -1 : link[x]
                      ]) {
                   tmp = match[y], match[x]
                         = y, match[y] = x;
                 return true;
               q.push(match[v]), vis[match[v
                    ]] = 1, dep[match[v]] = dep[u] + 2;
            } else if (vis[v] == 1 && find(
                  v) != find(u)) {
               int p = lca(u, v);
               blossom(u, v, p), blossom(v,
                    u, p);
         }
       }
       return false;
     int res = 0;
     for (int u = 0; u < n; ++u) { if (
           match[u] == -1) { res += augment}
           (u); } }
     return res:
  }
∫};
```

#### Kuhn Munkres

```
// need perfect matching or not : \ensuremath{\mathsf{w}}
     intialize with -INF / 0
template <typename Cost>
struct KM {
  static constexpr Cost INF =
       numeric_limits<Cost>::max() / 2;
  vector<Cost> hl, hr, slk;
  vector<int> l, r, pre, vl, vr;
  queue<int> q;
  vector<vector<Cost>> w;
  KM(int n) : n(n), hl(n), hr(n), slk(n),
        l(n, -1), r(n, -1), pre(n), vl(n)
        , vr(n),
    w(n, vector<Cost>(n, -INF)) {}
  bool check(int x) {
    vl[x] = true;
    if ([x] != -1) {
       q.push(l[x]);
       return vr[l[x]] = true;
    while (x != -1) \{ swap(x, r[l[x] =
         pre[x]]); }
    return false;
  void bfs(int s) {
    fill(slk.begin(), slk.end(), INF);
fill(vl.begin(), vl.end(), false);
fill(vr.begin(), vr.end(), false);
    q = \{\};
    q.push(s);
```

```
vr[s] = true;
  while (true) {
    Cost d;
    while (!q.empty()) {
      int y = q.front();
      q.pop();
      for (int x = 0; x < n; ++x) {
        if (!vl[x] \&\& slk[x] >= (d = hl)
             [x] + hr[y] - w[x][y])) {
          pre[x] = y;
          if (d != 0) {
            slk[x] = d;
          } else if (!check(x)) {
            return;
          }
        }
      }
    d = INF;
    for (int x = 0; x < n; ++x) { if (!
         vl[x] \&\& d > slk[x]) \{ d = slk \}
         [x]; }}
    for (int x = 0; x < n; ++x) {
      if (vl[x]) {
        hl[x] += d;
      } else {
        slk[x] -= d;
      if (vr[x]) { hr[x] -= d; }
    for (int x = 0; x < n; ++x) { if (!
         vl[x] && !slk[x] && !check(x))
          { return; }}
void addEdge(int u, int v, Cost x) { w[
    u][v] = max(w[u][v], x); }
Cost solve() {
  for (int i = 0; i < n; ++i) { hl[i] =</pre>
        *max_element(w[i].begin(), w[i
       ].end()); }
  for (int i = 0; i < n; ++i) { bfs(i);</pre>
  Cost res = 0;
  for (int i = 0; i < n; ++i) { res +=
      w[i][l[i]]; }
  return res;
```

### Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x, y, l, u), connect  $x \to y$  with capacity u - l.
  - For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v) > 0, connect  $S \to v$  with ca-If in(v) > 0, connect  $v \to 0$  pacity in(v), otherwise, connect  $v \to 0$  with capacity -in(v).
    - To maximize, connect  $t \rightarrow s$ with capacity  $\infty$  (skip this in circulation problem), and let fbe the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t
    - is the answer. To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f'\neq\sum_{v\in V,in(v)>0}in(v),$  there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e + f_e$ where  $f_e$  corresponds to the flow of edge e on the graph.
- · Construct minimum vertex cover from maximum matching M on bipartite graph
  - 1. Redirect every edge:  $y \rightarrow x$  if  $(x,y) \in M, x \rightarrow y$  otherwise. 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited.

- 4.  $y \in Y$  is chosen iff y is visited.
- Minimum cost cyclic flow
  - 1. Consruct super source S and sink T2. For each edge (x, y, c), connect  $x \rightarrow y$  with (cost, cap) = (c, 1) if c > y0, otherwise connect  $y \rightarrow x$  with (cost, cap) = (-c, 1)3. For each edge with c < 0, sum these
  - cost as K, then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v) > 0, connect  $S \to v$  with  $(\cos t, cap) =$ (0,d(v))
  - 5. For each vertex v with d(v) < 0, connect  $v \to T$  with  $(\cos t, cap) =$ (0, -d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer T
  - 2. Construct a max flow model, let *K* be the sum of all weights
  - 3. Connect source  $s \to v$ ,  $v \in G$  with

  - 3. Connect source s → v, v ∈ G with capacity K
    4. For each edge (u, v, w) in G, connect u → v and v → u with capacity w
    5. For v ∈ G, connect it with sink v → t with capacity K + 2T (∑<sub>e∈E(v)</sub> w(e)) 2w(v)
    6. T is a valid answer if the maximum
  - flow f < K|V|
- Minimum weight edge cover
  1. For each v ∈ V create a copy v', and connect u' → v' with weight w(u, v).
  2. Connect v → v' with weight 2µ(v),
  - where  $\mu(v)$  is the cost of the cheapest edge incident to v.

    3. Find the minimum weight perfect
  - matching on G'.
- Project selection problem
  - 1. If  $p_v > 0$ , create edge (s, v) with capacity  $p_v$ ; otherwise, create edge
  - (v,t) with capacity  $-p_v$ . 2. Create edge (u,v) with capacity w with w being the cost of choosing uwithout choosing v.
  - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming  $\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + xyx'y')$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity  $c_x$ and create edge (s, y) with capacity
- 2. Create edge (x, y) with capacity  $c_{xy}$ . 3. Create edge (x, y) and edge (x', y')with capacity  $c_{xyx'y'}$ .

## Data Structure

#### <ext/pbds> 3.1

// mergable heap

```
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef tree<int, null_type, std::less<</pre>
     int>, rb_tree_tag,
     tree_order_statistics_node_update>
     tree_set;
typedef cc_hash_table<int, int> umap;
typedef priority_queue<int> heap;
int main() {
  // rb tree
  tree_set s;
  s.insert(71); s.insert(22);
  assert(*s.find_by_order(0) == 22);
       assert(*s.find_by_order(1) == 71);
  assert(s.order_of_key(22) == 0); assert
       (s.order_of_key(71) == 1);
  s.erase(22);
  assert(*s.find_by_order(0) == 71);
       assert(s.order_of_key(71) == 0);
```

```
heap a, b; a.join(b);
   // persistant
  rope<char> r[2];
   r[1] = r[0];
  std::string st = "abc";
  r[1].insert(0, st.c_str());
r[1].erase(1, 1);
   std::cout << r[1].substr(0, 2) << std::
       endl;
  return 0:
| }
 3.2 Li Chao Tree
 constexpr i64 INF = 4e18;
struct Line {
   i64 a, b;
  Line(): a(0), b(INF) {}
   Line(i64 a, i64 b) : a(a), b(b) {}
   i64 \ operator()(i64 \ x) \ \{ \ return \ a \ * \ x +
 // [, ) !!!!!!!!!!
 struct Lichao {
  int n;
  vector<int> vals;
  vector<Line> lines;
  Lichao() {}
  void init(const vector<int> &v) {
    n = v.size();
    vals = v;
    sort(vals.begin(), vals.end());
    vals.erase(unique(vals.begin(), vals.
         end()), vals.end());
    lines.assign(4 * n, {});
  }
  int get(int x) { return lower_bound(
       vals.begin(), vals.end(), x)
       vals.begin(); }
   void apply(Line p, int id, int l, int r
       ) {
    Line &q = lines[id];
     if (p(vals[1]) < q(vals[1])) { swap(p
          , q); }
     if (l + 1 == r) { return; }
    int m = l + r \gg 1;
    if (p(vals[m]) < q(vals[m])) {</pre>
       swap(p, q);
       apply(p, id \ll 1, l, m);
       apply(p, id << 1 | 1, m, r);
    }
  void add(int ql, int qr, Line p) {
    ql = get(ql), qr = get(qr);
    auto go = [&](auto go, int id, int l,
           int r) -> void {
       if (ql <= l && r <= qr) {
         apply(p, id, l, r);
       int m = l + r >> 1;
       go(go, id << 1, 1, m);
       go(go, id << 1 | 1, m, r);
    go(go, 1, 0, n);
  i64 query(int p) {
    p = get(p);
    auto go = [&](auto go, int id, int l,
           int r) -> i64 {
       if (l + 1 == r) { return lines[id](
           vals[p]); }
       int m = l + r \gg 1;
       return min(lines[id](vals[p]), p <</pre>
            m ? go(go, id << 1, 1, m) : go
            (go, id << 1 | 1, m, r));
     return go(go, 1, 0, n);
∫};
      Treap
 3.3
struct Treap {
   Treap *lc = nullptr, *rc = nullptr;
```

```
int sz = 1;
    unsigned w = rng();
    i64 m = 0, b = 0, val = 0;
int size(Treap *t) {
    return t == nullptr ? 0 : t->sz;
void apply(Treap *t, i64 m, i64 b) {
    t->val += m * size(t->lc) + b;
void pull(Treap *t) {
    t\rightarrow sz = size(t\rightarrow lc) + size(t\rightarrow rc) +
void push(Treap *t) {
    if (t->lc != nullptr) {
        apply(t->lc, t->m, t->b);
    if (t->rc != nullptr) {
        apply(t->rc, t->m, t->b + t->m *
             (size(t->lc) + 1));
    t->m = t->b = 0;
pair<Treap*, Treap*> split(Treap *t, int
    s) {
    if (t == nullptr) { return {t, t}; }
    push(t);
    Treap *a, *b;
if (s <= size(t->lc)) {
   b = t;
        tie(a, b->lc) = split(t->lc, s);
    } else {
        a = t;
        tie(a->rc, b) = split(t->rc, s -
             size(t->lc) - 1);
    pull(t);
    return {a, b};
Treap* merge(Treap *t1, Treap *t2) {
    if (t1 == nullptr) { return t2; }
if (t2 == nullptr) { return t1; }
    push(t1), push(t2);
    if (t1->w > t2->w) {
        t1->rc = merge(t1->rc, t2);
        pull(t1);
        return t1;
    } else {
        t2->lc = merge(t1, t2->lc);
        pull(t2)
        return t2;
int rnk(Treap *t, i64 val) {
    int res = 0;
    while (t != nullptr) {
        push(t);
        if (val <= t->val) {
             res += size(t->lc) + 1;
             t = t->rc;
        } else {
             t = t->lc;
        }
    return res;
Treap* join(Treap *t1, Treap *t2) {
    if (size(t1) > size(t2)) {
        swap(t1, t2);
    Treap *t = nullptr;
    while (t1 != nullptr) {
        auto [u1, v1] = split(t1, 1);
        t1 = v1;
        int r = rnk(t2, u1->val);
        if (r > 0)
             auto [u2, v2] = split(t2, r);
             t = merge(t, u2);
             t2 = v2;
        t = merge(t, u1);
    t = merge(t, t2);
    return t;
```

```
3.4 Link-Cut Tree
struct Splay {
  array<Splay*, 2> ch = {nullptr, nullptr
  Splay* fa = nullptr;
  int sz = 1;
  bool rev = false;
  Splay() {}
  void applyRev(bool x) {
    if (x) {
      swap(ch[0], ch[1]);
rev ^= 1;
    }
  }
  void push() {
    for (auto k : ch) {
      if (k) {
        k->applyRev(rev);
      }
    rev = false;
  void pull() {
    sz = 1:
    for (auto k : ch) {
      if (k) {
    }
  int relation() { return this == fa->ch
  bool isRoot() { return !fa || fa->ch[0]
        != this && fa->ch[1] != this; }
  void rotate() {
    Splay *p = fa;
bool x = !relation();
    p \rightarrow ch[!x] = ch[x];
    if (ch[x]) { ch[x] \rightarrow fa = p; }
    fa = p -> fa:
    if (!p->isRoot()) { p->fa->ch[p->
         relation()] = this; }
    ch[x] = p;
    p \rightarrow fa = this;
    p->pull();
  void splay() {
    vector<Splay*> s;
    for (Splay *p = this; !p->isRoot(); p
          = p->fa) { s.push_back(p->fa);
         }
    while (!s.empty()) {
      s.back()->push();
      s.pop_back();
    }
    push();
while (!isRoot()) {
      if (!fa->isRoot()) {
        if (relation() == fa->relation())
          fa->rotate();
        } else {
          rotate();
        }
      }
      rotate();
    pull();
  void access() {
    for (Splay *p = this, *q = nullptr; p
         ; q = p, p = p -> fa) {
      p->splay();
      p->ch[1] = q;
      p->pull();
    splay();
  void makeRoot() {
    access();
    applyRev(true);
  Splay* findRoot() {
    access();
    Splay *p = this;
    while (p->ch[0]) \{ p = p->ch[0]; \}
    p->splay();
```

```
return p:
  friend void split(Splay *x, Splay *y) {
    x->makeRoot();
    y->access();
  }
  // link if not connected
  friend void link(Splay *x, Splay *y) {
    x->makeRoot();
    if (y->findRoot() != x) {
      x->fa=y;
  }
  // delete edge if doesn't exist
  friend void cut(Splay *x, Splay *y) {
    split(x, y);
    if (x->fa == y \&\& !x->ch[1]) {
       x->fa = y->ch[0] = nullptr;
       x->pull();
    }
  }
  bool connected(Splay *x, Splay *y) {
    return x->findRoot() == y->findRoot()
∫};
```

#### 4 Graph

## 2-Edge-Connected Components

```
struct EBCC {
  int n, cnt = 0, T = 0;
  vector<vector<int>> adj, comps;
  vector<int> stk, dfn, low, id;
  EBCC(int n) : n(n), adj(n), dfn(n, -1),
  low(n), id(n, -1) {}

void addEdge(int u, int v) { adj[u].
        push_back(v), adj[v].push_back(u);
  void build() { for (int i = 0; i < n; i
        ++) { if (dfn[i] == -1) { dfs(i,
        -1); }}}
  void dfs(int u, int p) {
    dfn[u] = low[u] = T++;
     stk.push_back(u);
     for (auto v : adj[u]) {
       if (v == p) { continue; }
if (dfn[v] == -1) {
         dfs(v, u);
         low[u] = min(low[u], low[v]);
       else\ if\ (id[v] == -1)
         low[u] = min(low[u], dfn[v]);
       }
     if (dfn[u] == low[u]) {
       int x;
       comps.emplace_back();
       do {
         x = stk.back();
         comps.back().push_back(x);
         id[x] = cnt;
         stk.pop_back();
       } while (x != u);
       cnt++;
  }
j};
```

#### 4.22-Vertex-Connected Components

```
// is articulation point if appear in >=
    2 comps
auto dfs = [&](auto dfs, int u, int p) ->
     void {
  dfn[u] = low[u] = T++;
  for (auto v : adj[u]) {
    if (v == p) { continue; }
    if (dfn[v] == -1) {
      stk.push_back(v);
      dfs(dfs, v, u);
      low[u] = min(low[u], low[v]);
```

```
if (low[v] >= dfn[u]) {
        comps.emplace_back();
        int x;
        do {
          x = stk.back();
          cnt[x]++;
          stk.pop_back();
        } while (x != v);
        comps.back().push_back(u);
        cnt[u]++;
   } else {
     low[u] = min(low[u], dfn[v]);
};
for (int i = 0; i < n; i++) {
 if (!adj[i].empty()) {
   dfs(dfs, i, -1);
 } else {
    comps.push_back({i});
```

## 3-Edge-Connected Components

```
// DSU
struct ETCC {
 int n, cnt = 0;
 vector<vector<int>> adj, comps;
  vector<int> in, out, low, up, nx, id;
 ETCC(int n): n(n), adj(n), in(n, -1), out(in), low(n), up(n), nx(in), id
       (in) {}
  void addEdge(int u, int v) {
    adj[u].push_back(v);
    adj[v].push_back(u);
 void build() {
    int T = 0;
    DSU d(n);
    auto merge = [&](int u, int v) {
      d.join(u, v);
      up[u] += up[v];
    auto dfs = [&](auto dfs, int u, int p
         ) -> void {
      in[u] = low[u] = T++;
      for (auto v : adj[u]) {
        if (v == u) { continue; }
        if (v == p) {
 p = -1;
          continue;
        if (in[v] == -1) {
          dfs(dfs, v, u);
          if (nx[v] == -1 \&\& up[v] <= 1)
            up[u] += up[v];
            low[u] = min(low[u], low[v]);
             continue;
          if (up[v] == 0) \{ v = nx[v]; \}
          if (low[u] > low[v]) \{ low[u] =
                 low[v], swap(nx[u], v); }
          while (v != -1) { merge(u, v);
                v = nx[v]; }
        } else if (in[v] < in[u]) {</pre>
          low[u] = min(low[u], in[v]);
          up[u]++;
        } else {
          for (int &x = nx[u]; x != -1 &&
                in[x] \leftarrow in[v] \& in[v] <
                out[x]; x = nx[x]) {
            merge(u, x);
          up[u]--;
        }
      }
      out[u] = T;
    for (int i = 0; i < n; i++) { if (in[</pre>
         i] == -1) { dfs(dfs, i, -1); }}
    for (int i = 0; i < n; i++) { if (d.
         find(i) == i) { id[i] = cnt++;
```

```
}}
     comps.resize(cnt);
     for (int i = 0; i < n; i++) { comps[</pre>
          id[d.find(i)]].push_back(i); }
  }
};
```

## Heavy-Light Decomposi-

```
struct HLD {
  int n, cur = 0;
  vector<int> sz, top, dep, par, tin,
       tout, seq;
  vector<vector<int>> adj;
  HLD(int n) : n(n), sz(n, 1), top(n),
       dep(n), par(n), tin(n), tout(n),
       seq(n), adj(n) {}
  void addEdge(int u, int v) { adj[u].
       push_back(v), adj[v].push_back(u);
  void build(int root = 0) {
    top[root] = root, dep[root] = 0, par[
         root7 = -1:
    dfs1(root), dfs2(root);
  void dfs1(int u) {
    if (auto it = find(adj[u].begin(),
         adj[u].end(), par[u]); it != adj
         [u].end()) {
      adj[u].erase(it);
    for (auto &v : adj[u]) {
      par[v] = u;
      dep[v] = dep[u] + 1;
      dfs1(v);
      sz[u] += sz[v];
      if (sz[v] > sz[adj[u][0]]) { swap(v
           , adj[u][0]); }
    }
  void dfs2(int u) {
    tin[u] = cur++:
    seq[tin[u]] = u;
    for (auto v : adj[u]) {
      top[v] = v == adj[u][0] ? top[u] :
      dfs2(v);
    tout[u] = cur - 1;
  int lca(int u, int v) {
    while (top[u] != top[v]) {
      if (dep[top[u]] > dep[top[v]]) {
        u = par[top[u]];
      } else {
        v = par[top[v]];
      }
    }
    return dep[u] < dep[v] ? u : v;</pre>
  int dist(int u, int v) { return dep[u]
       + dep[v] - 2 * dep[lca(u, v)]; }
  int jump(int u, int k) {
    if (dep[u] < k) { return -1; }
int d = dep[u] - k;</pre>
    while (dep[top[u]] > d) \{ u = par[top ] \}
         [u]]; }
    return seq[tin[u] - dep[u] + d];
  // u is v's ancestor
  bool isAncestor(int u, int v) { return
       tin[u] <= tin[v] && tin[v] <= tout
       [u]; }
  // root's parent is itself
  int rootedParent(int r, int u) {
    if (r == u) { return u; }
    if (isAncestor(r, u)) { return par[u
         ]; }
    auto it = upper_bound(adj[u].begin(),
          adj[u].end(), r, [\&](int x, int
      return tin[x] < tin[y];</pre>
```

}) - 1;

return \*it;

```
// rooted at u, v's subtree size
int rootedSize(int r, int u) {
  if (r == u) { return n; }
  if (isAncestor(u, r)) { return sz[u];
   }
  return n - sz[rootedParent(r, u)];
}
int rootedLca(int r, int a, int b) {
  return lca(a, b) ^ lca(a, r) ^ lca
  (b, r); }
};
```

## 4.5 Centroid Decomposition

```
vector<int> sz(n), vis(n);
auto build = [&](auto build, int u, int p
     ) -> void {
  sz[u] = 1;
  for (auto v : g[u]) {
    if (v != p && !vis[v]) {
      build(build, v, u);
      sz[u] += sz[v];
  }
};
auto find = [&](auto find, int u, int p,
     int tot) -> int {
  for (auto v : g[u]) {
    if (v != p && !vis[v] && 2 * sz[v] >
         tot) {
       return find(find, v, u, tot);
    }
  return u;
};
auto dfs = [%](auto dfs, int cen) -> void
  build(build, cen, -1);
  cen = find(find, cen, -1, sz[cen]);
  vis[cen] = 1;
  build(build, cen, -1);
  for (auto v : g[cen]) {
    if (!vis[v]) {
      dfs(dfs, v);
  }
}:
dfs(dfs, 0);
```

## 4.6 Strongly Connected Components

```
struct SCC {
  int n, cnt = 0, cur = 0;
  vector<int> id, dfn, low, stk;
  vector<vector<int>> adj, comps;
void addEdge(int u, int v) { adj[u].
       push_back(v); }
  SCC(int n): n(n), id(n, -1), dfn(n, -1), low(n, -1), adj(n) {}
  void build() {
    auto dfs = [&](auto dfs, int u) ->
         void {
      dfn[u] = low[u] = cur++;
      stk.push_back(u);
       for (auto v : adj[u]) {
         if (dfn[v] == -1) {
           dfs(dfs, v);
           low[u] = min(low[u], low[v]);
        } else if (id[v] == -1) {
           low[u] = min(low[u], dfn[v]);
        }
       if (dfn[u] == low[u]) {
         int v;
         comps.emplace_back();
         do {
           v = stk.back();
           comps.back().push_back(v);
           id[v] = cnt;
           stk.pop_back();
         } while (u != v);
         cnt++;
```

```
for (int i = 0; i < n; i++) { if (dfn
        [i] == -1) { dfs(dfs, i); }}
for (int i = 0; i < n; i++) { id[i] =
        cnt - 1 - id[i]; }
reverse(comps.begin(), comps.end());
}
// the comps are in topological sorted
order
};</pre>
```

#### 4.7 2-SAT

```
struct TwoSat {
   int n, N;
   vector<vector<int>> adj;
   vector<int> ans;
   TwoSat(int n) : n(n), N(n), adj(2 * n)
   // u == x
   void addClause(int u, bool x) { adj[2 *
   u + !x].push_back(2 * u + x); }
// u == x || v == y
   void addClause(int u, bool x, int v,
     bool y) {
adj[2 * u + !x].push_back(2 * v + y);
adj[2 * v + !y].push_back(2 * u + x);
   }
// u == x -> v == y
   void addImply(int u, bool x, int v,
         bool y) { addClause(u, !x, v, y);
   void addVar() {
      adj.emplace_back(), adj.emplace_back
           ();
   // at most one in var is true
   // adds prefix or as supplementary
         variables
   void atMostOne(const vector<pair<int,</pre>
         bool>> &vars) {
      int sz = vars.size();
     for (int i = 0; i < sz; i++) {
        addVar();
        auto [u, x] = vars[i];
        addImply(u, x, N - 1, true);
        if (i > 0) {
          addImply(N - 2, true, N - 1, true
          addClause(u, !x, N - 2, false);
     }
   // does not return supplementary
         variables from atMostOne()
   bool satisfiable() {
     // run tarjan scc on 2 * N \,
      for (int i = 0; i < 2 * N; i++) { if
     (dfn[i] == -1) { dfs(dfs, i); }}
for (int i = 0; i < N; i++) { if (id
    [2 * i] == id[2 * i + 1]) {</pre>
           return false; }}
      ans.resize(n);
      for (int i = 0; i < n; i++) { ans[i]
= id[2 * i] > id[2 * i + 1]; }
     return true;
   }
|};
```

## 4.8 count 3-cycles and 4-cycles

```
| sort(ord.begin(), ord.end(), [&](auto i, auto j) { return pair(deg[i], i) > pair(deg[j], j); }); | for (int i = 0; i < n; i++) { rnk[ord[i]] = i; } | if (rnk[u] < rnk[v]) { dag[u].push_back(v ); } | // c3 | for (int x = 0; x < n; x++) { for (auto y : dag[x]) { vis[y] = 1; } | for (auto y : dag[x]) { for (auto z : dag[y]) { ans += vis[z]; } | for (auto y : dag[x]) { vis[y] = 0; } | for (auto y : dag[x]) { vis[y] = 0; } | for (auto y : dag[x]) { vis[y] = 0; } | for (auto y : dag[x]) { vis[y] = 0; } | for (auto y : dag[x]) { vis[y] = 0; } | for (auto y : dag[x]) { vis[y] = 0; } | for (auto y : dag[x]) { vis[y] = 0; } | for (auto y : dag[x]) { vis[y] = 0; } | for (auto y : dag[x]) | for (auto y
```

```
| }
// c4
| for (int x = 0; x < n; x++) {
| for (auto y : dag[x]) { for (auto z :
| adj[y]) { if (rnk[z] > rnk[x]) {
| ans += vis[z]++; }}
| for (auto y : dag[x]) { for (auto z :
| adj[y]) { if (rnk[z] > rnk[x]) {
| vis[z]--; }}
| }
```

## 4.9 Minimum Mean Cycle

create a new vertex S, connect S to all vertices with arbitrary weight (0). Let  $f_i(u)$  be the shortest path from S to u with exactly i edges.

$$ans = \min_{f_{n+1}(i)! = \infty} \max_{j=1}^{n} \frac{f_{n+1}(i) - f_j(i)}{n+1-j}$$

## 4.10 Directed Minimum Spanning Tree

```
// DSU with rollback
template <typename Cost>
struct DMST {
 int n;
 vector<int> s, t, lc, rc, h;
vector<Cost> c, tag;
DMST(int n) : n(n), h(n, -1) {}
  void addEdge(int u, int v, Cost w) {
    int id = s.size();
    s.push_back(u), t.push_back(v), c.
         push_back(w);
    lc.push_back(-1), rc.push_back(-1);
    tag.emplace_back();
    h[v] = merge(h[v], id);
 pair<Cost, vector<int>> build(int root
       = 0) {
    DSU d(n);
    Cost res{};
    vector<int> vis(n, -1), path(n), q(n)
         , in(n, -1);
    vis[root] = root;
    vector<pair<int, vector<int>>> cycles
    for (auto r = 0; r < n; ++r) {
  auto u = r, b = 0, w = -1;
      while (!~vis[u]) {
        if (!~h[u]) { return {-1, {}}; }
        push(h[u]);
        int e = h[u];
        res += c[e], tag[h[u]] -= c[e];
        h[u] = pop(h[u]);
        q[b] = e, path[b++] = u, vis[u] =
        u = d.find(s[e]);
        if (vis[u] == r) {
          int cycle = -1, e = b;
            w = path[--b];
            cycle = merge(cycle, h[w]);
          } while (d.join(u, w));
          u = d.find(u);
          h[u] = cycle, vis[u] = -1;
          cycles.emplace_back(u, vector<
                int>(q.begin() + b, q.
               begin() + e);
        }
      for (auto i = 0; i < b; ++i) { in[d]
           .find(t[q[i]])] = q[i]; 
    reverse(cycles.begin(), cycles.end())
    for (const auto &[u, comp] : cycles)
      int count = int(comp.size()) - 1;
      d.back(count);
      int ine = in[u];
      for (auto e : comp) { in[d.find(t[e
           ])] = e; }
      in[d.find(t[ine])] = ine;
```

vector<int> par;

```
par.reserve(n);
     for (auto i : in) { par.push_back(i != -1 ? s[i] : -1); }
     return {res, par};
  void push(int u) {
     c[u] += tag[u];
     if (int l = lc[u]; l != -1) { tag[l]
          += tag[u]; }
     if (int r = rc[u]; r != -1) { tag[r]
         += tag[u]; }
     tag[u] = 0;
   int merge(int u, int v) {
     if (u == -1 || v == -1) { return u !=
          -1 ? u : v; }
     push(u);
     push(v);
     if (c[u] > c[v]) { swap(u, v); }
     rc[u] = merge(v, rc[u]);
     swap(lc[u], rc[u]);
     return u;
  int pop(int u) {
     push(u);
     return merge(lc[u], rc[u]);
|};
```

## 4.11 Maximum Clique

```
pair<int, vector<int>> maxClique(int n,
    const vector<bitset<N>> adj) {
  int mx = 0;
 vector<int> ans, cur;
 auto rec = [&](auto rec, bitset<N> s)
       -> void {
    int sz = s.count();
    if (int(cur.size()) > mx) { mx = cur.
         size(), ans = cur; }
    if (int(cur.size()) + sz <= mx) {</pre>
         return; }
    int e1 = -1, e2 = -1;
    vector<int> d(n);
for (int i = 0; i < n; i++) {</pre>
      if (s[i]) {
        d[i] = (adj[i] & s).count();
        if (e1 == -1 || d[i] > d[e1]) {
             e1 = i; }
        if (e2 == -1 || d[i] < d[e2]) {
             e2 = i; }
     }
    if (d[e1] >= sz - 2) {
      cur.push_back(e1);
      auto s1 = adj[e1] & s;
      rec(rec, s1);
      cur.pop_back();
      return;
    cur.push_back(e2);
    auto s2 = adj[e2] & s;
    rec(rec, s2);
    cur.pop_back();
    s.reset(e2);
    rec(rec, s);
 bitset<N> all;
 for (int i = 0; i < n; i++) {
   all.set(i);
 rec(rec, all);
  return pair(mx, ans);
```

## 4.12 Dominator Tree

```
// res : parent of each vertex in
   dominator tree, -1 is root, -2 if
   not in tree
| struct DominatorTree {
   int n, cur = 0;
   vector<int> dfn, rev, fa, sdom, dom,
      val, rp, res;
   vector<vector<int>> adj, rdom, r;
```

```
DominatorTree(int n) : n(n), dfn(n, -1)
         res(n, -2), adj(n), rdom(n), r(n)
       ) {
     rev = fa = sdom = dom = val = rp =
         dfn;
  }
   void addEdge(int u, int v) {
    adj[u].push_back(v);
   void dfs(int u) {
    dfn[u] = cur;
     rev[cur] = u;
     fa[cur] = sdom[cur] = val[cur] = cur;
     for (int v : adj[u]) {
       if (dfn[v] == -1) {
         dfs(v)
         rp[dfn[v]] = dfn[u];
       r[dfn[v]].push_back(dfn[u]);
  int find(int u, int c) {
     if (fa[u] == u) { return c != 0 ? -1
         : u; }
     int p = find(fa[u], 1);
     if (p == -1) { return c != 0 ? fa[u]
          : val[u]; }
     if (sdom[val[u]] > sdom[val[fa[u]]])
          { val[u] = val[fa[u]]; }
     fa[u] = p;
    return c != 0 ? p : val[u];
  void build(int s = 0) {
     dfs(s);
     for (int i = cur - 1; i >= 0; i--) {
       for (int u : r[i]) { sdom[i] = min(
            sdom[i], sdom[find(u, 0)]); }
       if (i > 0) { rdom[sdom[i]].
           push_back(i); }
       for (int u : rdom[i]) {
         int p = find(u, 0);
         if (sdom[p] == i) {
           dom[u] = i;
         } else {
           dom[u] = p;
       if (i > 0) { fa[i] = rp[i]; }
    res[s] = -1;
    for (int i = 1; i < cur; i++) { if (
         sdom[i] != dom[i]) { dom[i] =
         dom[dom[i]]; }}
     for (int i = 1; i < cur; i++) { res[</pre>
         rev[i]] = rev[dom[i]]; }
  }
|};
```

#### 4.13 Edge Coloring

```
// bipartite
e[i] = pair(u, v + a), deg[u]++, deg[v +
a]++;
int col = *max_element(deg.begin(), deg.
     end());
vector<int> ans(m, -1);
vector has(a + b, vector<pair<int, int>>(
col, \{-1, -1\}));
for (int i = 0; i < m; i++) {
  auto [u, v] = e[i];
  vector<int> c;
  for (auto x : {u, v}) {
    c.push_back(0);
    while (has[x][c.back()].first != -1)
          { c.back()++; }
  if (c[0] != c[1]) {
    auto dfs = [\&](auto dfs, int u, int x
         ) -> void {
      auto [v, i] = has[u][c[x]];
      if (v != -1) {
         if (has[v][c[x ^ 1]].first != -1)
           dfs(dfs, v, x ^ 1);
        } else {
```

```
has[v][c[x]] = \{-1, -1\};
        has[u][c[x ^ 1]] = \{v, i\}, has[v]
             ][c[x \land 1]] = \{u, i\};
        ans[i] = c[x \wedge 1];
      }
    dfs(dfs, v, 0);
  has[u][c[0]] = \{v, i\};
  has[v][c[0]] = {u, i};
  ans[i] = c[0];
// general
auto vizing(int n, const vector<pair<int,</pre>
      int>> &e) {
  vector<int> deg(n);
for (auto [u, v] : e) {
    deg[u]++, deg[v]++;
  int col = *max_element(deg.begin(), deg
       .end()) + 1;
  vector<int> free(n);
  vector ans(n, vector<int>(n, -1));
  vector at(n, vector<int>(col, -1));
  auto update = [&](int u) {
    free[u] = 0;
    while (at[u][free[u]] != -1) {
      free[u]++;
    }
  };
  auto color = [&](int u, int v, int c1)
    int c2 = ans[u][v];
    ans[u][v] = ans[v][u] = c1;
    at[u][c1] = v, at[v][c1] = u;
    if (c2 != -1) {
      at[u][c2] = at[v][c2] = -1;
      free[u] = free[v] = c2;
    } else {
      update(u), update(v);
    return c2;
 };
  auto flip = [&](int u, int c1, int c2)
    int v = at[u][c1];
    swap(at[u][c1], at[u][c2]);
    if (v != -1) {
      ans[u][v] = ans[v][u] = c2;
    if (at[u][c1] == -1) {
      free[u] = c1;
    if (at[u][c2] == -1) {
      free[u] = c2;
    return v:
  for (int i = 0; i < int(e.size()); i++)</pre>
    auto [u, v1] = e[i];
    int v2 = v1, c1 = free[u], c2 = c1, d
    vector<pair<int, int>> fan;
    vector<int> vis(col);
while (ans[u][v1] == -1) {
      fan.emplace_back(v2, d = free[v2]);
      if (at[v2][c2] == -1) {
        for (int j = int(fan.size()) - 1;
              j >= 0; j--) {
          c2 = color(u, fan[j].first, c2)
      } else if (at[u][d] == -1) {
        for (int j = int(fan.size()) - 1;
               j >= 0; j--) +
          color(u, fan[j].first, fan[j].
               second);
      } else if (vis[d] == 1) {
        break;
      } else {
        vis[d] = 1, v2 = at[u][d];
      }
    if (ans[u][v1] == -1) {
```

## 5 String

#### 5.1 Prefix Function

```
template <typename T>
vector<int> prefixFunction(const T &s) {
  int n = int(s.size());
  vector<int> p(n);
  for (int i = 1; i < n; i++) {
    int j = p[i - 1];
    while (j > 0 && s[i] != s[j]) { j = p
        [j - 1]; }
    if (s[i] == s[j]) { j++; }
    p[i] = j;
  }
  return p;
}
```

## 5.2 Z Function

## 5.3 Suffix Array

```
// need to discretize
struct SuffixArray {
  int n:
vector<int> sa, as, ha;
template <typename T>
  vector<int> sais(const T &s) {
     int n = s.size(), m = *max_element(s.
          begin(), s.end()) + 1;
     vector < int > pos(m + 1), f(n);
    for (auto ch : s) { pos[ch + 1]++; } for (int i = 0; i < m; i++) { pos[i +
           1] += pos[i]; }
     for (int i = n - 2; i >= 0; i--) { f[
          i] = s[i] != s[i + 1] ? s[i] < s
          [i + 1]: f[i + 1]; }
    vector<int> x(m), sa(n);
    auto induce = [&](const vector<int> &
          ls) {
       fill(sa.begin(), sa.end(), -1);
auto L = [&](int i) { if (i >= 0 &&
              !f[i]) { sa[x[s[i]]++] = i;}
             }};
```

```
auto S = [\&](int i) \{ if (i >= 0 \&\&
             f[i] { sa[--x[s[i]]] = i;
            }};
       for (int i = 0; i < m; i++) { x[i]</pre>
            = pos[i + 1]; }
       for (int i = int(ls.size()) - 1; i
            >= 0; i--) { S(ls[i]); }
       for (int i = 0; i < m; i++) { x[i]
            = pos[i]; }
       L(n - 1);
       for (int i = 0; i < n; i++) { L(sa[</pre>
            i] - 1); }
       S(sa[i] - 1); }
     auto ok = [&](int i) { return i == n
     || !f[i - 1] && f[i]; };
auto same = [&](int i, int j) {
       do { if (s[i++] != s[j++]) { return
             false; }} while (!ok(i) && !
            ok(j));
       return ok(i) && ok(j);
     vector<int> val(n), lms;
     for (int i = 1; i < n; i++) { if (ok(
          i)) { lms.push_back(i); }}
     induce(lms);
     if (!lms.empty()) {
       int p = -1, w = 0;
for (auto v : sa) {
         if (v != 0 && ok(v)) {
           if (p != -1 \&\& same(p, v)) \{ w \}
                --; }
           val[p = v] = w++;
         }
       auto b = lms;
       for (auto &v : b) { v = val[v]; }
       b = sais(b);
       for (auto &v : b) { v = lms[v]; }
       induce(b);
     return sa;
 template <typename T>
   SuffixArray(const T &s) : n(s.size()),
       sa(sais(s)), as(n), ha(n - 1) {
     for (int i = 0; i < n; i++) { as[sa[i
         ]] = i; }
     for (int i = 0, j = 0; i < n; ++i) {
       if (as[i] == 0) {
         j = 0;
       } else {
         for (j -= j > 0; i + j < n & sa[
              as[i] - 1] + j < n && s[i +
              j] == s[sa[as[i] - 1] + j];
                { ++j; }
         ha[as[i] - 1] = j;
};
```

## 5.4 Manacher's Algorithm

## 5.5 Aho-Corasick Automaton

```
constexpr int K = 26;
struct Node {
  array<int, K> nxt;
  int fail = -1;
// other vars
  Node() { nxt.fill(-1); }
};
vector<Node> aho(1);
for (int i = 0; i < n; i++) {</pre>
  string s;
  cin >> s;
  int u = 0;
  for (auto ch : s) {
    int c = ch - 'a'
    if (aho[u].nxt[c] == -1) {
      aho[u].nxt[c] = aho.size();
      aho.emplace_back();
    u = aho[u].nxt[c];
 }
vector<int> q;
for (auto &i : aho[0].nxt) {
  if (i == -1) {
    i = 0;
  } else {
    q.push_back(i);
    aho[i].fail = 0;
  }
for (int i = 0; i < int(q.size()); i++) {</pre>
  int u = q[i];
  if (u > 0) {
    // maintain
  for (int c = 0; c < K; c++) {</pre>
    if (int v = aho[u].nxt[c]; v != -1) {
      aho[v].fail = aho[aho[u].fail].nxt[
           c];
      q.push_back(v);
      aho[u].nxt[c] = aho[aho[u].fail].
           nxt[c];
 }
```

#### 5.6 Suffix Automaton

```
struct SAM {
  static constexpr int A = 26;
  struct Node {
    int len = 0, link = -1, cnt = 0;
    array<int, A> nxt;
    Node() { nxt.fill(-1); }
  vector<Node> t;
  SAM() : t(1) {}
  int size() { return t.size(); }
  Node& operator[](int i) { return t[i];
  int newNode() {
    t.emplace_back();
    return t.size() - 1;
  int extend(int p, int c) {
    int cur = newNode();
    t[cur].len = t[p].len + 1;
    t[cur].cnt = 1;
    while (p != -1 && t[p].nxt[c] == -1)
      t[p].nxt[c] = cur;
      p = t[p].link;
    if (p == -1) {
      t[cur].link = 0;
    } else {
      int q = t[p].nxt[c];
      if (t[p].len + 1 == t[q].len) {
        t[cur].link = q;
      } else {
        int clone = newNode();
        t[clone].len = t[p].len + 1;
```

## 5.7 Lexicographically Smallest Rotation

```
template <typename T>
T minRotation(T s) {
 int n = s.size();
 int i = 0, j = 1;
s.insert(s.end(), s.begin(), s.end());
 while (i < n && j < n) {</pre>
    int k = 0;
    while (k < n \&\& s[i + k] == s[j + k])
    if (s[i + k] <= s[j + k]) {
      j += k + 1;
    } else {
      i += k + 1;
    if (i == j) {
      j++;
  int ans = i < n ? i : j;
 return T(s.begin() + ans, s.begin() +
       ans + n;
```

## 5.8 EER Tree

```
// cnt : occurrences, (dfs fail tree)
// num : number of pal ending here
struct PAM {
  static constexpr int A = 26;
  struct Node {
    int len = 0, link = 0, cnt = 0, num =
         0;
    array<int, A> nxt{};
    Node() {}
  vector<Node> t;
  int suf = 1;
  string s;
  PAM() : t(2) { t[0].len = -1; }
  int size() { return t.size(); }
  Node& operator[](int i) { return t[i];
  int newNode() {
    t.emplace_back();
    return t.size() - 1;
  bool add(int c, char offset = 'a') {
    int pos = s.size();
    s += c + offset;
    int cur = suf, curlen = 0;
    while (true) {
      curlen = t[cur].len;
      if (pos - 1 - curlen >= 0 && s[pos
           - 1 - curlen] == s[pos]) {
           break; }
      cur = t[cur].link;
    if (t[cur].nxt[c]) {
      suf = t[cur].nxt[c];
      t[suf].cnt++;
      return false;
    suf = newNode();
    t[suf].len = t[cur].len + 2;
    t[suf].cnt = t[suf].num = 1;
    t[cur].nxt[c] = suf;
    if (t[suf].len == 1) {
```

## 6 Math

#### 6.1 Extended GCD

```
| array<i64, 3> extgcd(i64 a, i64 b) {
| if (b == 0) { return {a, 1, 0}; }
| auto [g, x, y] = extgcd(b, a % b);
| return {g, y, x - a / b * y};
| }
```

## 6.2 Chinese Remainder Theorem

```
// returns (rem, mod), n = 0 return (0,
     1), no solution return (0, 0)
pair<i64, i64> crt(vector<i64> r, vector<
     i64 > m) {
  int n = r.size();
  for (int i = 0; i < n; i++) {
    r[i] %= m[i];
    if (r[i] < 0) { r[i] += m[i]; }</pre>
  i64 \ r0 = 0, \ m0 = 1;
  for (int i = 0; i < n; i++) {
    i64 r1 = r[i], m1 = m[i];
    if (m0 < m1) { swap(r0, r1), swap(m0,</pre>
          m1); }
    if (m0 \% m1 == 0) {
      if (r0 % m1 != r1) { return {0, 0};
      continue:
    auto [g, a, b] = extgcd(m0, m1);
    i64 u1 = m1 / g;
    if ((r1 - r0) % g != 0) { return {0,
         0}; }
    i64 x = (r1 - r0) / g % u1 * a % u1; r0 += x * m0; m0 *= u1;
    if (r0 < 0) \{ r0 += m0; \}
  return {r0, m0};
```

#### 6.3 NTT and polynomials

```
template <int P>
struct Modint {
  int v;
 // need constexpr, constructor, +-*,
      qpow, inv()
template<int P>
constexpr Modint<P> findPrimitiveRoot() {
 Modint < P > i = 2;
  int k = __builtin_ctz(P - 1);
 while (true) {
    if (i.qpow((P - 1) / 2).v != 1) {
        break; }
   i = i + 1:
 }
 return i.qpow(P - 1 >> k);
template <int P>
constexpr Modint<P> primitiveRoot =
    findPrimitiveRoot<P>();
vector<int> rev;
```

template <int P>

```
vector<Modint<P>> roots{0, 1};
template <int P>
void dft(vector<Modint<P>> &a) {
  int n = a.size();
  if (n == 1) { return; }
  if (int(rev.size()) != n) {
    int k = __builtin_ctz(n) - 1;
    rev.resize(n);
    for (int i = 0; i < n; i++) { rev[i]</pre>
         = rev[i >> 1] >> 1 | (i & 1) <<
  for (int i = 0; i < n; i++) { if (rev[i
       ] < i) { swap(a[i], a[rev[i]]); }}
  if (roots<P>.size() < n) {</pre>
    int k = __builtin_ctz(roots<P>.size()
         );
    roots<P>.resize(n);
    while ((1 << k) < n) {
      auto e = Modint<P>(primitiveRoot<P</pre>
           >).qpow(P - 1 >> k + 1);
      for (int i = 1 \ll k - 1; i < 1 \ll k
           ; i++) {
        roots<P>[2 * i] = roots<P>[i];
roots<P>[2 * i + 1] = roots<P>[i]
      k++:
    }
  }
  for (int k = 1; k < n; k *= 2) {
  for (int i = 0; i < n; i += 2 * k) {
    for (int j = 0; j < k; j++) {
      Modint<P> u = a[i + j];
    }
}
        Modint < P > v = a[i + j + k] *
              roots<P>[k + j];
        // fft : v = a[i + j + k] * roots

[n / (2 * k) * j]
         a[i + j] = u + v;
        a[i + j + k] = u - v;
      }
    }
 }
template <int P>
void idft(vector<Modint<P>> &a) {
  int n = a.size();
  reverse(a.begin() + 1, a.end());
  dft(a);
  Modint < P > x = (1 - P) / n;
  for (int i = 0; i < n; i++) { a[i] = a[
       `i] * x; }
template <int P>
struct Poly : vector<Modint<P>>> {
  using Mint = Modint<P>;
  Poly() {}
  explicit Poly(int n) : vector<Mint>(n)
      {}
  explicit Poly(const vector<Mint> &a) :
       vector<Mint>(a) {}
  explicit Poly(const initializer_list<</pre>
Mint> &a) : vector<Mint>(a) {} template<class F>
  explicit Poly(int n, F f) : vector<Mint</pre>
       >(n) { for (int i = 0; i < n; i++) }
        { (*this)[i] = f(i); }}
template<class InputIt>
  explicit constexpr Poly(InputIt first,
       InputIt last) : vector<Mint>(first
        , last) {}
  Poly mulxk(int k) {
    auto b = *this;
    b.insert(b.begin(), k, 0);
    return b;
  Poly modxk(int k) {
    k = min(k, int(this->size()));
    return Poly(this->begin(), this->
         begin() + k);
  Poly divxk(int k) {
    if (this->size() <= k) { return Poly</pre>
```

```
return Poly(this->begin() + k, this->
        end());
friend Poly operator+(const Poly &a,
     const Poly &b) {
  Poly res(max(a.size(), b.size()));
  for (int i = 0; i < int(a.size()); i
       ++) { res[i] = res[i] + a[i]; }
  for (int i = 0; i < int(b.size()); i</pre>
       ++) { res[i] = res[i] + b[i]; }
  return res;
Poly res(max(a.size(), b.size()));
  for (int i = 0; i < int(a.size()); i</pre>
  ++) { res[i] = res[i] + a[i]; }
for (int i = 0; i < int(b.size()); i
++) { res[i] = res[i] - b[i]; }
  return res;
friend Poly operator*(Poly a, Poly b) {
  if (a.empty() || b.empty()) { return
       Poly(); }
  int sz = 1, tot = a.size() + b.size()
        - 1;
  while (sz < tot) { sz *= 2; }
  a.resize(sz);
  b.resize(sz);
  dft(a);
  dft(b);
  for (int i = 0; i < sz; i++) { a[i] =
    a[i] * b[i]; }</pre>
  idft(a);
  a.resize(tot);
  return a;
friend Poly operator*(Poly a, Mint b) {
  for (int i = 0; i < int(a.size()); i
++) { a[i] = a[i] * b; }
  return a;
Poly derivative() {
  if (this->empty()) { return Poly(); }
  Poly res(this->size() - 1);
  for (int i = 0; i < this->size() - 1;
++i) { res[i] = (i + 1) * (*
       this)[i + 1]; }
  return res;
Poly integral() {
  Poly res(this->size() + 1);
  for (int i = 0; i < this->size(); ++i
) { res[i + 1] = (*this)[i] *
       Mint(i + 1).inv(); }
  return res;
Poly inv(int m) {
  // a[0] != 0
  Poly x({(*this)[0].inv()});
  int k = 1;
  while (k < m) {</pre>
    k *= 2;
x = (x * (Poly({2}) - modxk(k) * x)
          ).modxk(k);
  return x.modxk(m);
Poly log(int m) {
  return (derivative() * inv(m)).
       integral().modxk(m);
Poly exp(int m) {
  Poly x(\{1\});
  int k = 1;
  while (k < m) {
   k *= 2;
   x = (x * (Poly({1}) - x.log(k) +</pre>
          modxk(k)).modxk(k);
  return x.modxk(m);
Poly pow(i64 k, int m) {
  if (k == 0) { return Poly(m, [&](int
        i) { return i == 0; }); }
  int i = 0;
```

```
while (i < this->size() && (*this)[i
    ].v == 0) { i++; }
if (i == this->size() || ...
                                   int128(i)
          * k >= m) { return Poly(m); }
    Mint v = (*this)[i];
    auto f = divxk(i) * v.inv();
return (f.log(m - i * k) * k).exp(m -
           i * k).mulxk(i * k) * v.qpow(k)
  Poly sqrt(int m) {
    // a[0] == 1, otherwise quadratic
          residue?
    Poly x(\{1\});
    int k = 1;
    while (k < m) {
   k *= 2;</pre>
       x = (x + (modxk(k) * x.inv(k)).
            modxk(k)) * ((P + 1) / 2);
    return x.modxk(m);
  Poly mulT(Poly b) const {
    if (b.empty()) { return Poly(); }
    int n = b.size();
    reverse(b.begin(), b.end());
    return (*this * b).divxk(n - 1);
  vector<Mint> evaluate(vector<Mint> x) {
    if (this->empty()) { return vector<</pre>
          Mint>(x.size()); }
    int n = max(x.size(), this->size());
    vector<Poly> q(4 * n);
    vector<Mint> ans(x.size());
    x.resize(n);
    auto build = [&](auto build, int id,
          int l, int r) -> void {
       if (r - l == 1) {
         q[id] = Poly(\{1, -x[l].v\});
       } else {
         int m = (l + r) / 2;
build(build, 2 * id, l, m);
build(build, 2 * id + 1, m, r);
q[id] = q[2 * id] * q[2 * id +
               1];
      }
    build(build, 1, 0, n);
    auto work = [&](auto work, int id,
          int 1, int r, const Poly &num)
          -> void {
       if (r - l == 1) {
         if (l < int(ans.size())) { ans[l]</pre>
                = num[0]; }
      } else {
         int m = (1 + r) / 2;
work(work, 2 * id, 1, m, num.mulT
        (q[2 * id + 1]).modxk(m - 1)
         }
    work(work, 1, 0, n, mulT(q[1].inv(n))
    return ans;
template <int P>
Poly<P> interpolate(vector<Modint<P>> x,
     vector<Modint<P>> y) {
  // f(xi) = yi
  int n = x.size();
  vector<Poly<P>> p(4 * n), q(4 * n);
  auto dfs1 = [&](auto dfs1, int id, int
    l, int r) -> void {
if (l == r) {
       p[id] = Poly < P > (\{-x[l].v, 1\});
    int m = l + r >> 1;
    fit m = 1 + 1 >> 1,
dfs1(dfs1, id << 1, 1, m);
dfs1(dfs1, id << 1 | 1, m + 1, r);
p[id] = p[id << 1] * p[id << 1 | 1];</pre>
```

```
dfs1(dfs1, 1, 0, n - 1);
Poly<P> f = Poly<P>(p[1].derivative().
           evaluate(x));
    auto dfs2 = [&](auto dfs2, int id, int
           1, int r) -> void {
       if (1 == r) {
          q[id] = Poly < P > (\{y[l] * f[l].inv()
                 });
          return;
       int m = l + r >> 1;
      dfs2(dfs2, id << 1, 1, m);
dfs2(dfs2, id << 1 | 1, m + 1, r);
q[id] = q[id << 1] * p[id << 1 | 1] +
    q[id << 1 | 1] * p[id << 1];</pre>
    dfs2(dfs2, 1, 0, n - 1);
    return q[1];
 auto shift = [&](FPS f, int k) {
    FPS a(n + 1), b(n + 1);
    Mint powk = 1;
   for (int i = 0; i <= n; i++) {
    a[i] = ifact[i] * powk;
    b[i] = fact[i] * f[i];
    powk = powk * k;
    reverse(b.begin(), b.end());
auto g = a * b;
    g.resize(n + 1);
    reverse(g.begin(), g.end());
    for (int i = 0; i <= n; i++) {
  g[i] = g[i] * ifact[i];</pre>
    return g;
};
```

## 6.4 Any Mod NTT

```
constexpr int P0 = 998244353, P1 =
    1004535809, P2 = 469762049;
constexpr i64 P01 = 1LL * P0 * P1;
constexpr int inv0 = Modint<P1>(P0).inv()
    .v;
constexpr int inv01 = Modint<P2>(P01).inv
    ().v;
for (int i = 0; i < int(c.size()); i++) {
    i64 x = 1LL * (c1[i] - c0[i] + P1) % P1
        * inv0 % P1 * P0 + c0[i];
    c[i] = ((c2[i] - x % P2 + P2) % P2 *
        inv01 % P2 * (P01 % P) % P + x) %
    P;
}</pre>
```

## 6.5 Newton's Method

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

## 6.6 Fast Walsh-Hadamard Transform

```
1. XOR Convolution
```

```
• f(A) = (f(A_0) + f(A_1), f(A_0) - f(A_1))
• f^{-1}(A) = (f^{-1}(\frac{A_0 + A_1}{2}), f^{-1}(\frac{A_0 - A_1}{2}))
```

2. OR Convolution

• 
$$f(A) = (f(A_0), f(A_0) + f(A_1))$$
  
•  $f^{-1}(A) = (f^{-1}(A_0), f^{-1}(A_1) - f^{-1}(A_0))$ 

3. AND Convolution

• 
$$f(A) = (f(A_0) + f(A_1), f(A_1))$$
  
•  $f^{-1}(A) = (f^{-1}(A_0) - f^{-1}(A_1), f^{-1}(A_1))$ 

## 6.7 Simplex Algorithm

Description: maximize  $\mathbf{c}^T\mathbf{x}$  subject to  $A\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Returns  $-\infty$  if infeasible and  $\infty$  if unbounded.

```
const double eps = 1e-9;
 const double inf = 1e+9;
 int n, m;
 vector<double>> d;
 vector<int> p, q;
 void pivot(int r, int s) {
   double inv = 1.0 / d[r][s];
   for (int i = 0; i < m + 2; ++i) {
  for (int j = 0; j < n + 2; ++j) {
    if (i != r && j != s) d[i][j] -= d[
            r][j] * d[i][s] * inv;
     }
   }
   for (int i = 0; i < m + 2; ++i) if (i
        != r) d[i][s] *= -inv;
   for (int j = 0; j < n + 2; ++j) if (j
        != s) d[r][j] *= +inv;
   d[r][s] = inv;
   swap(p[r], q[s]);
 bool phase(int z) {
   int x = m + z;
   while (true) {
     int s = -1;
     for (int i = 0; i <= n; ++i) {
       if (!z && q[i] == -1) continue;
       if (s == -1] d[x][i] < d[x][s]) s
     if (d[x][s] > -eps) return true;
     int r = -1;
     for (int i = 0; i < m; ++i) {
       if (d[i][s] < eps) continue;
if (r == -1 || d[i][n + 1] / d[i][s</pre>
            ] < d[r][n + 1] / d[r][s]) r =
     if (r == -1) return false;
     pivot(r, s);
  }
 }
 vector<double> solve(const vector<vector<
      double>> &a, const vector<double> &b
      , const vector<double> &c) {
   m = b.size(), n = c.size();
   d = vector<vector<double>>(m + 2,
        vector<double>(n + 2));
   for (int i = 0; i < m; ++i) {
  for (int j = 0; j < n; ++j) d[i][j] =</pre>
           a[i][j];
   p.resize(m), q.resize(n + 1);
   for (int i = 0; i < m; ++i) p[i] = n +
        i, d[i][n] = -1, d[i][n + 1] = b[i
        ];
   for (int i = 0; i < n; ++i) q[i] = i, d
        [m][i] = -c[i];
   q[n] = -1, d[m + 1][n] = 1;
   int r = 0;
   for (int i = 1; i < m; ++i) if (d[i][n + 1] < d[r][n + 1]) r = i;
   if (d[r][n + 1] < -eps) {
     pivot(r, n);
     if (!phase(1) || d[m + 1][n + 1] < -
           eps) return vector<double>(n, -
          inf);
     for (int i = 0; i < m; ++i) if (p[i]
          == -1) {
        int s = min_element(d[i].begin(), d
             [i].end() - 1) - d[i].begin();
       pivot(i, s);
   }
   if (!phase(0)) return vector<double>(n,
         inf);
   vector<double> x(n);
   for (int i = 0; i < m; ++i) if (p[i] <</pre>
        n) x[p[i]] = d[i][n + 1];
   return x;
| }
```

#### 6.7.1 Construction

Standard form: maximize  $\mathbf{c}^T \mathbf{x}$  subject to  $A\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Dual LP: minimize  $\mathbf{b}^T \mathbf{y}$  subject to  $A^T \mathbf{y} \geq \mathbf{c}$  and  $\mathbf{y} \geq 0$ .

```
\bar{\mathbf{x}} and \bar{\mathbf{y}} are optimal if and only if for all i \in [1, n], either \bar{x}_i = 0 or \sum_{j=1}^m A_{ji}\bar{y}_j = c_i holds and for all i \in [1, m] either \bar{y}_i = 0 or \sum_{j=1}^m A_{ij}\bar{x}_j = b_j holds.
```

- $\begin{array}{ll} \text{1. In case of minimization, let } c_i' = -c_i \\ \text{2. } \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq \\ -b_j \\ \text{3. } \sum_{1 \leq i \leq n} A_{ji} x_i = b_j \end{array}$ 
  - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$   $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$   $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x'_i$

### 6.8 Subset Convolution

```
Description: h(s) = \sum_{s' \subset s} f(s')g(s \setminus s')
```

```
vector<int> SubsetConv(int n, const
     vector<int> &f, const vector<int> &g
      ) {
  const int m = 1 \ll n;
  vector<vector<int>> a(n + 1, vector<int
  >(m)), b(n + 1, vector<int>(m));
for (int i = 0; i < m; ++i) {
     a[__builtin_popcount(i)][i] = f[i];
     b[__builtin_popcount(i)][i] = g[i];
  for (int i = 0; i <= n; ++i) {
  for (int j = 0; j < n; ++j) {</pre>
       for (int s = 0; s < m; ++s) {
          if (s >> j & 1) {
            a[i][s] += a[i][s ^ (1 << j)];
b[i][s] += b[i][s ^ (1 << j)];
       }
    }
  vector<vector<int>> c(n + 1, vector<int
        >(m));
  for (int s = 0; s < m; ++s) {
    for (int i = 0; i <= n; ++i) {
  for (int j = 0; j <= i; ++j) c[i][s
   ] += a[j][s] * b[i - j][s];
    }
  for (int i = 0; i <= n; ++i) {</pre>
     for (int j = 0; j < n; ++j) {
       for (int s = 0; s < m; ++s) {
  if (s >> j & 1) c[i][s] -= c[i][s
                 ^ (1 << j)];
    }
  }
  vector<int> res(m);
  for (int i = 0; i < m; ++i) res[i] = c[</pre>
          _builtin_popcount(i)][i];
  return res;
```

## $\begin{array}{ccc} \textbf{6.9} & \textbf{Berlekamp Massey Algorithm} \\ \end{array}$

```
// find \sum a_(i-j)c_j = 0 for d \le i template <typename T>
vector<T> berlekampMassey(const vector<T>
      &a) {
  vector<T> c(1, 1), oldC(1);
  int oldI = -1;
  T \text{ oldD} = 1;
  for (int i = 0; i < int(a.size()); i++)</pre>
    T d = 0;
    for (int j = 0; j < int(c.size()); j</pre>
          ++) { d += c[j] * a[i - j]; }
    if (d == 0) { continue; }
T mul = d / oldD;
    vector<T> nc = c;
    nc.resize(max(int(c.size()), i - oldI
           + int(oldC.size()));
    for (int j = 0; j < int(oldC.size());</pre>
          j++) { nc[j + i - oldI] -= oldC
[j] * mul; }
    if (i - int(c.size()) > oldI - int(
          oldC.size())) {
       oldI = i;
       oldD = d:
```

```
swap(oldC, c);
}
swap(c, nc);
}
return c;
}
```

#### 6.10 Fast Linear Recurrence

```
// p : a[0] \sim a[d - 1]
// q : a[i] = \sum a[i - j]q[j]
template <typename T>
T linearRecurrence(vector<T> p, vector<T>
       q, i64 n) {
   int d = q.size() - 1;
  assert(int(p.size()) == d);
p = p * q;
   p.resize(d);
   while (n > 0) {
     auto nq = q;
     for (int i = 1; i <= d; i += 2) {
       nq[i] *= -1;
     auto np = p * nq;
nq = q * nq;
     for (int i = 0; i < d; i++) {
  p[i] = np[i * 2 + n % 2];
     for (int i = 0; i <= d; i++) {
   q[i] = nq[i * 2];
     n /= 2:
  }
   return p[0] / q[0];
```

## 6.11 Prime check and factorize

```
i64 mul(i64 a, i64 b, i64 mod) {}
i64 qpow(i64 x, i64 p, i64 mod) {}
bool isPrime(i64 n) {
  if (n == 1) { return false; }
  int r = __builtin_ctzll(n - 1);
  i64 d = n - 1 >> r;
  auto checkComposite = [&](i64 p) {
    i64 x = qpow(p, d, n);
     if (x == 1 \mid \mid x == n - 1) { return
          false; }
     for (int i = 1; i < r; i++) {
       x = mul(x, x, n);
       if (x == n - 1) { return false; }
    return true;
  for (auto p : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}) {
     if (n == p) {
  return true;
    } else if (checkComposite(p)) {
       return false;
  return true;
vector<i64> pollardRho(i64 n) {
  vector<i64> res;
  auto work = [&](auto work, i64 n) {
     if (n <= 10000) {
      for (int i = 2; i * i <= n; i++) {
  while (n % i == 0) {
           res.push_back(i);
           n /= i;
       if (n > 1) { res.push_back(n); }
       return;
     } else if (isPrime(n)) {
       res.push_back(n);
       return;
     i64 \times 0 = 2;
    auto f = [\&](i64 x) \{ return (mul(x, extension)) \}
         x, n) + 1) % n; };
     while (true) {
       i64 x = x0, y = x0, d = 1, power = 1, lam = 0, v = 1;
```

```
while (d == 1) {
      y = f(y);
++lam;
       v = mul(v, abs(x - y), n);
       if (lam % 127 == 0) {
         d = gcd(v, n);
v = 1;
       if (power == lam) {
         x = y;
power *= 2;
         lam = 0;
         d = gcd(v, n);
v = 1;
    if (d != n) {
      work(work, d);
       work(work, n / d);
       return;
    ++x0;
  }
};
work(work, n);
sort(res.begin(), res.end());
return res;
```

### 6.12 Count Primes leq n

```
// __attribute__((target("avx2"),
optimize("03", "unroll-loops")))
i64 primeCount(const i64 n) {
   if (n <= 1) { return 0; }
if (n == 2) { return 1; }</pre>
   const int v = sqrtl(n);
   int s = (v + 1) / 2;
   vector<int> smalls(s), roughs(s), skip(
        v + 1);
   vector<i64> larges(s);
   iota(smalls.begin(), smalls.end(), 0);
   for (int i = 0; i < s; i++) {
  roughs[i] = 2 * i + 1;</pre>
     larges[i] = (n / roughs[i] - 1) / 2;
   const auto half = [](int n) -> int {
        return (n - 1) >> 1; };
   int pc = 0;
   for (int p = 3; p \leftarrow v; p += 2) {
     if (skip[p]) { continue; }
     int q = p * p;
if (1LL * q * q > n) { break; }
     skip[p] = true;
      for (int i = q; i \le v; i += 2 * p)
           skip[i] = true;
      int ns = 0;
      for (int k = 0; k < s; k++) {
        int i = roughs[k];
        if (skip[i]) { continue; }
i64 d = 1LL * i * p;
        larges[ns] = larges[k] - (d \ll v)
              larges[smalls[d / 2] - pc] :
              smalls[half(n / d)]) + pc;
        roughs[ns++] = i;
      s = ns;
     for (int i = half(v), j = v / p - 1 |
1; j >= p; j -= 2) {
int c = smalls[j / 2] - pc;
for (int e = j * p / 2; i >= e; i
              --) { smalls[i] -= c; }
     pc++;
   larges[0] += 1LL * (s + 2 * (pc - 1)) *
   (s - 1) / 2;
for (int k = 1; k < s; k++) { larges[0]
          -= larges[k]; }
   for (int l = 1; l < s; l++) {
     i64 q = roughs[1];
     i64 M = n / q;
int e = smalls[half(M / q)] - pc;
      if (e <= 1) { break; }</pre>
      i64 t = 0;
      for (int k = l + 1; k \le e; k++) { t
           += smalls[half(M / roughs[k])];
```

## 6.13 Discrete Logarithm

```
| // return min x >= 0 s.t. a ^ x = b mod m
       , 0 \land 0 = 1, -1 if no solution
 // (I think) if you want x > 0 (m != 1),
      remove if (b == k) return add;
 int discreteLog(int a, int b, int m) {
   if (m == 1) {
     return 0;
   a %= m, b %= m;
   int k = 1, add = 0, g;
   while ((g = gcd(a, m)) > 1) {
  if (b == k) {
        return add;
      } else if (b % g) {
        return -1;
     b /= g, m /= g, ++add;
k = 1LL * k * a / g % m;
   if (b == k) {
     return add;
   int n = sqrt(m) + 1;
   int an = 1;
for (int i = 0; i < n; ++i) {
   an = 1LL * an * a % m;</pre>
   unordered_map<int, int> vals;
   for (int q = 0, cur = b; q < n; ++q) {
     vals[cur] = q;
cur = 1LL * a * cur % m;
   for (int p = 1, cur = k; p <= n; ++p) {
  cur = 1LL * cur * an % m;</pre>
      if (vals.count(cur)) {
        int ans = n * p - vals[cur] + add;
        return ans;
     }
   return -1;
```

#### 6.14 Quadratic Residue

```
1// rna
 int jacobi(int a, int m) {
   int s = 1;
   while (m > 1) {
     a %= m;
     if (a == 0) { return 0; }
     int r = __builtin_ctz(a);
if (r % 2 == 1 && (m + 2 & 4) != 0) {
           s = -s; 
     a \gg = r:
     if ((a \& m \& 2) != 0) \{ s = -s; \}
     swap(a, m);
   return s;
 int quadraticResidue(int a, int p) {
   if (p == 2) { return a % 2; }
   int j = jacobi(a, p);
   if (j == 0 \mid | j == -1) \{ return j; \}
   int b, d;
   while (true) {
     b = rng() % p;
d = (1LL * b * b + p - a) % p;
     if (jacobi(d, p) == -1) \{ break; \}
   int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp
   for (int e = p + 1 >> 1; e > 0; e >>=
        1) {
     if (e % 2 == 1) {
       tmp = (1LL * g0 * f0 + 1LL * d * g1
       % p * f1 % p) % p;
g1 = (1LL * g0 * f1 + 1LL * g1 * f0
            ) % p;
       g0 = tmp;
```

## 6.15 Characteristic Polynomial

```
vector<vector<int>> Hessenberg(const
     vector<vector<int>> &A) {
  int N = A.size();
  vector<vector<int>> H = A;
  for (int i = 0; i < N - 2; ++i) {
  if (!H[i + 1][i]) {
       for (int j = i + 2; j < N; ++j) {
         if (H[j][i]) {
           for (int k = i; k < N; ++k)
                 swap(H[i + 1][k], H[j][k])
           for (int k = 0; k < N; ++k)
                 swap(H[k][i + 1], H[k][j])
           break;
         }
      }
    if (!H[i + 1][i]) continue;
int val = fpow(H[i + 1][i], kP - 2);
    for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP
      for (int k = i; k < N; ++k) H[j][k]
= (H[j][k] + 1LL * H[i + 1][k]
      ] * (kP - coef)) % kP;
for (int k = 0; k < N; ++k) H[k][i
            + 1] = (H[k][i + 1] + 1LL * H[
            k][j] * coef) % kP;
   }
  }
  return H;
}
vector<int> CharacteristicPoly(const
     vector<vector<int>> &A) {
  int N = A.size();
  auto H = Hessenberg(A);
  for (int i = 0; i < N; ++i) {
  for (int j = 0; j < N; ++j) H[i][j] =
           kP - H[i][j];
  vector<vector<int>>> P(N + 1, vector<int</pre>
       >(N + 1));
  P[0][0] = 1;
  for (int i = 1; i \le N; ++i) {
    P[i][0] = 0;
    for (int j = 1; j <= i; ++j) P[i][j]</pre>
          = P[i - 1][j - 1];
    int val = 1;
    for (int j = i - 1; j >= 0; --j) {
      int coef = 1LL * val * H[j][i - 1]
      coef) % kP;
if (j) val = 1LL * val * (kP - H[j
            ][j - 1]) % kP;
    }
  if (N & 1) {
    for (int i = 0; i <= N; ++i) P[N][i]</pre>
          = kP - P[N][i];
  return P[N];
```

## 6.16 Linear Sieve Related

```
mobius[i] = -1;
     for (int p : primes) {
         if (p > N / i) {
   break;
        minp[p * i] = p;
mobius[p * i] = -mobius[i];
if (i % p == 0) {
    mobius[p * i] = 0;
    head;
             break;
     }
| }
```

## De Bruijn Sequence

```
int res[kN], aux[kN], a[kN], sz;
void Rec(int t, int p, int n, int k) {
  if (t > n) {
     if (n \% p == 0)
       for (int i = 1; i <= p; ++i) res[sz
             ++] = aux[i];
    aux[t] = aux[t - p];
Rec(t + 1, p, n, k);
for (aux[t] = aux[t - p] + 1; aux[t]
           < k; ++aux[t]) Rec(t + 1, t, n,
  }
int DeBruijn(int k, int n) {
  // return cyclic string of length k^n
        such that every string of length n
         using k character appears as a
        substring.
  if (k == 1) return res[0] = 0, 1;
fill(aux, aux + k * n, 0);
return sz = 0, Rec(1, 1, n, k), sz;
```

#### Floor Sum 6.18

```
| // \sum \{i = 0\} \{n\} floor((a * i + b) / c
 i64 floorSum(i64 a, i64 b, i64 c, i64 n)
    if (n < 0) { return 0; }</pre>
    if (n == 0) { return b / c; }
if (a == 0) { return b / c * (n + 1); }
    i64 \text{ res} = 0;
    if (a >= c) { res += a / c * n * (n +
1) / 2, a %= c; }
if (b >= c) { res += b / c * (n + 1), b
    %= c; }
i64 m = (a * n + b) / c;
    return res + n * m - (m == 0 ? 0 :
floorSum(c, c - b - 1, a, m - 1));
|}
```

## 6.19 More Floor Sum

•  $m = \lfloor \frac{an+b}{c} \rfloor$ 

$$g(a,b,c,n) = \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \qquad \qquad (i,j) \text{ in } G.$$

$$= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} & - \text{ The num in } G \text{ is } | d \text{ in } G \text{ is } | d \text{ in } G \text{ is } | d \text{ in } G \text{ is } | d \text{ in } G \text{ is } | d \text{ in } G \text{ is } | d \text{ in } G \text{ is } | d \text{ in } G \text{ is } | d \text{ in } G \text{ is } | d \text{ in } G \text{ is } | d \text{ in } G \text{ is } | d \text{ in } G \text{ in } G \text{ in } | d \text{ in } G \text{ in } G \text{ in } | d \text{ in } G \text{ in }$$

$$h(a,b,c,n) = \sum_{i=0}^{n} \left\lfloor \frac{ai+b}{c} \right\rfloor^{2} \qquad \qquad \text{and } \frac{(x,y) \in E, \text{ our } \frac{aih(D)}{2} \text{ is the maxin}}{\frac{rank(D)}{2}} \text{ is the maxin}$$

$$= \begin{cases} \left\lfloor \frac{a}{c} \right\rfloor^{2} \cdot \frac{n(n+1)(2n+1)}{2} + \left\lfloor \frac{b}{c} \right\rfloor^{2} \cdot (n+1) \\ + \left\lfloor \frac{a}{c} \right\rfloor \cdot \left\lfloor \frac{b}{c} \right\rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \left\lfloor \frac{a}{c} \right\rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \left\lfloor \frac{b}{c} \right\rfloor \cdot f(a \bmod c, b \bmod c, c, n), \qquad a \geq c \lor b \geq c \\ 0, \qquad n < 0 \lor a \neq 0 \end{pmatrix} \qquad \text{tices, there are}$$

$$= \begin{cases} nm(m+1) - 2g(c, c-b-1, a, m-1) \\ -2f(c, c-b-1, a, m-1) - f(a, b, c, n), \qquad \text{otherwise spanning trees.} \end{cases}$$

#### 6.20Min Mod Linear

```
|// \min i : [0, n) (a * i + b) % m
// ok in 1e9
int minModLinear(int n, int m, int a, int
        b, int cnt = 1, int p = 1, int q =
       1) {
   if (a == 0) { return b; } if (cnt % 2 == 1) {
      if (b >= a) {
        int t = (m - b + a - 1) / a;
int c = (t - 1) * p + q;
        if (n <= c) { return b; }
n -= c;
b += a * t - m;
     }
b = a - 1 - b;
   } else {
      if (b < m - a) {
        int t = (m - b - 1) / a;
        int c = t * p;
if (n <= c) { return (n - 1) / p *
        a + b; }
n -= c;
b += a * t;
     b = m - 1 - b;
   cnt++;
   int d = m / a;
   int c = minModLinear(n, a, m % a, b,
   cnt, (d - 1) * p + q, d * p + q);
return cnt % 2 == 1 ? m - 1 - c : a - 1
           - c;
```

## Count of subsets with sum (mod P) leq T

```
int n, T;
 cin >> n >> T;
 vector<int> cnt(T + 1);
 for (int i = 0; i < n; i++) {
  int a;
  cin >> a;
  cnt[a]++;
 vector<Mint> inv(T + 1);
 for (int i = 1; i <= T; i++) {
  inv[i] = i == 1 ? 1 : -P / i * inv[P %
 FPS f(T + 1);
f = f.exp(T + 1);
```

#### 6.22Theorem

• Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning  $a \geq c \ \forall \text{three rooted at } r \text{ in } G \text{ is } |\det(\tilde{L}_{rr})|.$

oLet where  $d_{ij} = x_{ij}$  oLet where  $d_{ij} = x_{ij}$  of  $(x_{ij} \text{ is chosen uniformly at random})$  if i < j and  $(i, j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

- Cayley's Formula
  - Given Given a degree sequence  $d_1, d_2, \ldots, d_n$  for each labeled vertices, there are

$$\begin{array}{l} a \geq c \vee b \geq c & (n-2)! \\ n < 0 \vee a \ \overline{\not=d-1)!(d_2-1)! \cdots (d_n-1)!} \end{array}$$

- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1, 2, \dots, k$  belong to different components. Then  $T_{n,k} = kn^{n-k-1}$ .

#### • Erdős-Gallai Theorem

A sequence of non-negative integers  $d_1 \ge$  $d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1 + d_2 + \ldots + d_n$ 

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all  $1 \leq k \leq n$ .

• Burnside's Lemma

Let X be a set and G be a group that acts on X. For  $g \in G$ , denote by  $\hat{X}^g$  the elements fixed by g:

$$X^g = \{x \in X \mid gx \in X\}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

• Gale–Ryser theorem

A pair of sequences of nonnegative integers If pair of sequences of inclinations are integers  $a_1 \ge \cdots \ge a_n$  and  $b_1, \ldots, b_n$  is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \le$ 

 $\sum_{i=1}^{n} \min(b_i, k) \text{ holds for every } 1 \leq k \leq n.$ 

Sequences a and b called bigraphic if there is a labeled simple bipartite graph such that a and b is the degree sequence of this bipartite graph.

 $\bullet \quad \text{Fulkerson--Chen--Anstee theorem} \\$ 

A sequence  $(a_1, b_1), \ldots, (a_n, b_n)$  of nonnegative integer pairs with  $a_1 \geq \cdots \geq a_n$  $a_n$  is digraphic if and only if  $\sum_{i=1}^{n} a_i =$ 

$$\sum_{i=1}^{n} b_{i} \text{ and } \sum_{i=1}^{k} a_{i} \leq \sum_{i=1}^{k} \min(b_{i}, k-1) + \sum_{i=1}^{n} a_{i} \leq \sum_{i=1}^{k} a_{i} \leq \sum_{i=1}^{n} a$$

 $\sum_{=k+1}^{n} \min(b_i, k) \text{ holds for every } 1 \leq k \leq n.$ 

Sequences a and b called digraphic if there is a labeled simple directed graph such that each vertex  $v_i$  has indegree  $a_i$  and outdegree  $b_i$ .

• Pick's theorem

For simple polygon, when are all integer, we have  $\#\{\text{lattice points in the interior}\}\ \frac{\#\{\text{lattice points on the boundary}\}}{2}-1$ when points

• Möbius inversion formula

$$\begin{array}{lll} - & f(n) &= \sum_{d\mid n} g(d) & \Leftrightarrow & g(n) &= \\ & \sum_{d\mid n} \mu(d) f(\frac{n}{d}) \\ - & f(n) &= \sum_{n\mid d} g(d) & \Leftrightarrow & g(n) &= \\ & \sum_{n\mid d} \mu(\frac{d}{n}) f(d) \end{array}$$

- · Spherical cap
  - A portion of a sphere cut off by a
  - plane.
    r: sphere radius, a: radius of the base of the cap, h: height of the cap,
  - Example 3 and cap, n. height of the cap,  $\theta$ :  $\arcsin(a/r)$ , V volume =  $\pi h^2(3r-h)/3 = \pi h(3a^2+h^2)/6 = \pi r^3(2+\cos\theta)(1-\cos\theta)^2/3$ . Area =  $2\pi rh$  =  $\pi(a^2+h^2)$  =  $2\pi r^2(1-\cos\theta)$ .
- The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200000 for n < 1e19.
- The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30 for  $n=0\sim 9,$  627 for n=20,  $\sim 2e5$  for  $n=50, \sim 2e8$  for n=100.

```
• Total number of partitions of 
 n distinct elements: B(n) = 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, 115975, 
 <math>27644437, 190899322, \dots
```

- Ordinary Generating Function  $A(x) = \sum_{i \geq 0} a_i x^i$ 

```
-A(rx) \Rightarrow r^n a_n
-A(x) + B(x) \Rightarrow a_n + b_n
-A(x)B(x) \Rightarrow \sum_{i=0}^n a_i b_{n-i}
-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}
-xA(x)' \Rightarrow na_n
-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^n a_i
```

- Exponential Generating Function  $A(x) = \sum_{i \geq 0} \frac{a_i}{i!} x_i$ 

```
\begin{array}{ll} - & A(x) + B(x) \Rightarrow a_n + b_n \\ - & A^{(k)}(x) \Rightarrow a_{n \pm k_n} \\ - & A(x)B(x) \Rightarrow \sum_{i=0}^{k_n} \binom{n}{i} a_i b_{n-i} \\ - & A(x)^k \Rightarrow \sum_{i_1 + i_2 + \dots + i_k = n} \binom{n}{(i_1, i_2, \dots, i_k)} \\ - & xA(x) \Rightarrow na_n \end{array}
```

• Special Generating Function

$$- (1+x)^{n} = \sum_{i \ge 0} {n \choose i} x^{i}$$

$$- \frac{1}{(1-x)^{n}} = \sum_{i \ge 0} {n \choose i-1} x^{i}$$

• Bernoulli numbers

$$B_0 = 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, \text{ EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k^+ n^{m+1-k}$$

Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

```
\begin{array}{lll} S(n,k) &=& S(n-1,k-1) \, + \, kS(n-1,k), S(n,1) = S(n,n) = 1 \\ S(n,k) &=& \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n &=& \sum_{i=0}^n S(n,i)(x)_i \end{array}
```

• Pentagonal number theorem

$$\begin{array}{ll} \prod_{n=1}^{\infty}(1 \ -x^n) &= 1 \\ \sum_{k=1}^{\infty}(-1)^k\left(x^{k(3k+1)/2} + x^{k(3k-1)/2}\right) \end{array} +$$

• Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ , k+1 j:s s.t.  $\pi(j) \geq j$ , k j:s s.t.  $\pi(j) > j$ . E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k) E(n,0) = E(n,n-1) = 1  $E(n,k) = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n$ 

# Dynamic Programming

## 7.1 Dynamic Convex Hull

```
bool isect(iterator x, iterator y) {
     if (y == end()) return x->p = INF, 0;
if (x->k == y->k) x->p = x->b > y->b
? INF : -INF;
      else x->p = div(y->b - x->b, x->k - y
           ->k);
     return x->p >= y->p;
   void add(i64 k, i64 b) {
     auto z = insert(\{k, b, 0\}), y = z++,
           x = y;
     while (isect(y, z)) z = erase(z);
if (x != begin() && isect(--x, y))
           isect(x, y = erase(y));
     while ((y = x) != begin() \&\& (--x)->p
            >= V->p)
        isect(x, erase(y));
   i64 query(i64 x) {
     if (empty()) {
       return -INF;
     auto l = *lower_bound(x);
     return l.k * x + l.b;
|};
```

## 7.2 1D/1D Convex Optimization

```
struct segment {
  int i, l, r;
  segment(int a, int b, int c): i(a), l(b
       ), r(c) {}
inline long long f(int l, int r) { return
     dp[l] + w(l + 1, r); }
void solve() {
  dp[0] = 011;
  deque<segment> deq; deq.push_back(
  segment(0, 1, n));
for (int i = 1; i <= n; ++i) {
  dp[i] = f(deq.front().i, i);
  while (deq.size() && deq.front().r < i</pre>
       + 1) deq.pop_front();
  deq.front().l = i + 1;
  segment seg = segment(i, i + 1, n);
  while (deq.size() && f(i, deq.back().1)
        < f(deq.back().i, deq.back().l))
       deq.pop_back();
  if (deq.size()) {
    int d = 1048576, c = deq.back().1;
    while (d \gg 1) if (c + d \ll deq.back)
         ().r) {
    if(f(i, c + d) > f(deq.back().i, c +
          d)) c += d;
    deq.back().r = c; seg.l = c + 1;
  if (seg.l <= n) deq.push_back(seg);</pre>
```

## 7.3 Condition

## 7.3.1 Totally Monotone (Concave/Convex)

```
\begin{array}{ll} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \implies B[i][j'] \leq \\ B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \implies B[i][j'] \geq \\ B[i'][j'] \end{array}
```

## 7.3.2 Monge Condition (Concave/ Convex)

```
 \begin{array}{l} \forall i < i', j < j', B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

#### 7.3.3 Optimal Split Point

```
B[i][j] + B[i+1][j+1] \ge B[i][j+1] + B[i+1][j]
```

```
then H_{i,j-1} \le H_{i,j} \le H_{i+1,j}
```

## 8 Ckisemetry

### 8.1 Basic

```
#define IM imag
#define RE real
using lld = int64_t;
using llf = long double;
using PT = complex<lld>;
using PF = complex<llf>;
using P = PT;
llf abs(P p) { return sqrtl(norm(p)); }
PF toPF(PT p) { return PF{RE(p), IM(p)};
int sgn(lld x) \{ return (x > 0) - (x < 0) \}
lld dot(P a, P b) { return RE(conj(a) * b
    ); }
lld cross(P a, P b) { return IM(conj(a) *
b); }
int ori(P a, P b, P c) {
 return sgn(cross(b - a, c - a));
int quad(P p) {
  return (IM(p) == \emptyset) // use sgn for PF
    ? (RE(p) < 0 ? 3 : 1) : (IM(p) < 0 ?
int argCmp(P a, P b) {
  // returns 0/+-1, starts from theta = -
  int qa = quad(a), qb = quad(b);
  if (qa != qb) return sgn(qa - qb);
  return sgn(cross(b, a));
P rot90(P p) { return P{-IM(p), RE(p)}; }
template <typename V> llf area(const V &
    pt) {
  11d ret = 0; // BE CAREFUL OF TYPE!
  for (int i = 1; i + 1 < (int)pt.size();</pre>
    ret += cross(pt[i] - pt[0], pt[i+1] -
 pt[0]);
return ret / 2.0;
template <typename V> PF center(const V &
  pt) {
P ret = 0; lld A = 0; // BE CAREFUL OF
       TYPFI
  for (int i = 1; i + 1 < (int)pt.size();</pre>
        i++) {
    lld cur = cross(pt[i] - pt[0], pt[i
         +1] - pt[0]);
    ret += (pt[i] + pt[i + 1] + pt[0]) * cur; A += cur;
 }
  return toPF(ret) / llf(A * 3);
PF project(PF p, PF q) { // p onto q
  return dot(p, q) * q / dot(q, q); //
       dot<llf>
```

#### 8.2 ConvexHull

### 8.3 CyclicTS

```
int cyclic_ternary_search(int N, auto &&
      lt_) {
   auto lt = [\&](int x, int y) {
     return lt_(x % N, y % N); };
  int l = 0, r = N; bool up = lt(0, 1);
while (r - l > 1) {
  int m = (l + r) / 2;
     if (lt(m, 0) ? up : !lt(m, m+1)) r =
     else l' = m:
  }
   return (lt(l, r) ? r : l) % N;
} // find maximum; be careful if N == 0
```

## 8.4 Delaunay

```
/* please ensure input points are unique
/* A triangulation such that no points
     will strictly
inside circumcircle of any triangle. C
     should be big
enough s.t. the initial triangle contains
      all points */
#define L(i) ((i)==0 ? 2 : (i)-1)
#define R(i) ((i)==2 ? 0 : (i)+1)
#define F3 for (int i = 0; i < 3; i++)
bool is_inf(P z) { return RE(z) <= -C ||</pre>
     RE(z) >= C; }
bool in_cc(const array<P,3> &p, P q) {
  i128 inf_det = 0, det = 0, inf_N, N;
  F3 {
    if (is_inf(p[i]) && is_inf(q))
         continue;
    else if (is_inf(p[i])) inf_N = 1, N =
           -norm(q);
    else if (is_inf(q)) inf_N = -1, N =
         norm(p[i]);
    else inf_N = \emptyset, N = norm(p[i]) - norm
         (q);
    lld D = cross(p[R(i)] - q, p[L(i)] -
         q);
    inf_det += inf_N * D; det += N * D;
  return inf_det != 0 ? inf_det > 0 : det
P v[maxn];
struct Tri;
struct E {
  Tri *t; int side;
E(Tri *t_=0, int side_=0) : t(t_), side
       (side_) {}
struct Tri {
  array<int,3> p; array<Tri*,3> ch; array
     <E,3> e;
  Tri(int a=0, int b=0, int c=0) : p{a, b , c}, ch{} {}
  bool has_chd() const { return ch[0] !=
       nullptr; }
  bool contains(int q) const {
    F3 if (ori(v[p[i]], v[p[R(i)]], v[q])
          < 0)
       return false;
    return true;
  bool check(int q) const {
    return in_cc({v[p[0]], v[p[1]], v[p
[2]]}, v[q]); }
} pool[maxn * 10], *it, *root;
/* SPLIT_HASH_HERE */
void link(const E &a, const E &b) {
  if (a.t) a.t->e[a.side] = b;
  if (b.t) b.t->e[b.side] = a;
void flip(Tri *A, int a) {
  auto [B, b] = A->e[a]; /* flip edge
    between A,B */
  if (!B || !A->check(B->p[b])) return;
  Tri *X = new (it++) Tri(A->p[R(a)], B->
       p[b], A->p[a]);
  Tri *Y = new (it++) Tri(B->p[R(b)], A->
       p[a], B->p[b]);
  link(E(X, 0), E(Y, 0));
```

```
link(E(X, 1), A\rightarrow e[L(a)]); link(E(X, 2))
         B \rightarrow e[R(b)];
 link(E(Y, 1), B\rightarrow e[L(b)]); link(E(Y, 2))
       , A->e[R(a)]);
  A \rightarrow ch = B \rightarrow ch = \{X, Y, nullptr\};
 flip(X, 1); flip(X, 2); flip(Y, 1);
flip(Y, 2);
void add_point(int p) {
 Tri *r = root;
 while (r->has_chd()) for (Tri *c: r->ch
    if (c && c->contains(p)) { r = c;
         break; }
 array<Tri*, 3> t; /* split into 3
       triangles */
  F3 t[i] = new (it++) Tri(r->p[i], r->p[
 R(i)], p);
F3 link(E(t[i], 0), E(t[R(i)], 1));
  F3 link(E(t[i], 2), r->e[L(i)]);
 F3 flip(t[i], 2);
auto build(const vector<P> &p) {
 it = pool; int n = (int)p.size();
  vector<int> ord(n); iota(all(ord), 0);
 shuffle(all(ord), mt19937(114514));
 root = new (it++) Tri(n, n + 1, n + 2);
  copy_n(p.data(), n, v); v[n++] = P(-C,
       -C);
  v[n++] = P(C * 2, -C); v[n++] = P(-C, C)
        * 2);
 for (int i : ord) add_point(i);
 vector<array<int, 3>> res;
  for (Tri *now = pool; now != it; now++)
    if (!now->has_chd()) res.push_back(
         now->p);
 return res;
```

### 8.5 DirInPoly

```
| bool DIP(const auto &p, int i, P dir) {
  const int n = (int)p.size();
  P A = p[i+1==n ? 0 : i+1] - p[i];
  P B = p[i==0 ? n-1 : i-1] - p[i];
  if (auto C = cross(A, B); C < 0)
    return cross(A, dir) >= 0 || cross(
         dir, B) >= 0;
  else
    return cross(A, dir) >= 0 && cross(
         dir, B) >= 0;
} // is Seg(p[i], p[i]+dir*eps) in p? (
     non-strict)
// p is counterclockwise simple polygon
```

#### 8.6 FarthestPair

```
// p is CCW convex hull w/o colinear
     points
int n = (int)p.size(), pos = 1; lld ans =
   0;
for (int i = 0; i < n; i++) {
  P = p[(i + 1) \% n] - p[i];
  while (cross(e, p[(pos + 1) % n] - p[i
        ]) >
          cross(e, p[pos] - p[i]))
     pos = (pos + 1) % n;
   for (int j: {i, (i + 1) % n})
ans = max(ans, norm(p[pos] - p[j]));
} // tested @ AOJ CGL_4_B
```

## 8.7 HPIGeneralLine

```
struct Line {
  lld a, b, c; // ax + by + c <= 0
P dir() const { return P(a, b); }
  Line(lld ta, lld tb, lld tc) : a(ta), b
         (tb), c(tc) {}
  Line(P S, P T):a(IM(T-S)),b(-RE(T-S)),c
(cross(T,S)) {}
}; using LN = const Line &;
PF intersect(LN A, LN B) {
  llf c = cross(A.dir(), B.dir());
i128 a = i128(A.c) * B.a - i128(B.c) *
         A.a;
```

```
i128 b = i128(A.c) * B.b - i128(B.c) *
       A.b:
  return PF(-b / c, a / c);
bool cov(LN 1, LN A, LN B) {
  i128 c = cross(A.dir(), B.dir());
  i128 \ a = i128(A.c) * B.a - i128(B.c) *
       A.a;
  i128 b = i128(A.c) * B.b - i128(B.c) *
       A.b;
  return sgn(a * l.b - b * l.a + c * l.c)
        * sgn(c) >= 0;
bool operator<(LN a, LN b) {</pre>
  if (int c = argCmp(a.dir(), b.dir()))
       return c == -1
  return i128(abs(b.a) + abs(b.b)) * a.c
                   i128(abs(a.a) + abs(a.b)
                        )) * b.c;
|}
```

#### 8.8 HalfPlaneIntersection

```
struct Line {
  P st, ed, dir;
  Line (P s, P e) : st(s), ed(e), dir(e -
        s) {}
}; using LN = const Line &;
PF intersect(LN A, LN B) {
  llf t = cross(B.st - A.st, B.dir) /
    llf(cross(A.dir, B.dir));
  return toPF(A.st) + toPF(A.dir) * t; //
        C^3 / C^2
bool cov(LN 1, LN A, LN B) {
  i128 u = cross(B.st-A.st, B.dir);
  i128 v = cross(A.dir, B.dir);
  // ori(l.st, l.ed, A.st + A.dir*(u/v))
        <= 0?
   i128 \times = RE(A.dir) * u + RE(A.st - l.st
        ) * v;
  i128 y = IM(A.dir) * u + IM(A.st - l.st
) * v;
  return sgn(x*IM(l.dir) - y*RE(l.dir)) *
        sgn(v) >= 0;
} // x, y are C^3, also sgn<i128> is needed
bool operator<(LN a, LN b) {</pre>
  if (int c = argCmp(a.dir, b.dir))
    return c == -1:
  return ori(a.st, a.ed, b.st) < 0;</pre>
// cross(pt-line.st, line.dir)<=0 <-> pt
      in half plane
// the half plane is the LHS when going
      from st to ed
llf HPI(vector<Line> &q) {
  sort(q.begin(), q.end());
   int n = (int)q.size(), l = 0, r = -1;
  for (int i = 0; i < n; i++) {
     if (i && !argCmp(q[i].dir, q[i-1].dir
          )) continue;
     while (l < r && cov(q[i], q[r-1], q[r
         ])) --r;
     while (l < r \&\& cov(q[i], q[l], q[l]
          +1])) ++1;
     q[++r] = q[i];
  while (l < r \&\& cov(q[l], q[r-1], q[r])
       ) --r;
   while (l < r && cov(q[r], q[l], q[l+1])</pre>
        ) ++1;
  n = r - l + 1; // q[l .. r] are the
        lines
   if (n \le 2 \mid | argCmp(q[l].dir, q[r].
        dir)) return 0;
  vector<PF> pt(n);
for (int i = 0; i < n; i++)</pre>
    pt[i] = intersect(q[i+l], q[(i+1)%n+l])
          ]);
  return area(pt);
} // test @ 2020 Nordic NCPC : BigBrother
```

#### 8.9 HullCut

#### **8.10** KDTree

```
struct KDTree {
   struct Node {
     int x, y, x1, y1, x2, y2, id, f; Node *L, *R;
   } tree[maxn], *root;
   lld dis2(int x1, int y1, int x2, int y2
     lld dx = x1 - x2, dy = y1 - y2;
return dx * dx + dy * dy;
   static bool cmpx(Node& a, Node& b) {
        return a.x<b.x: }
   static bool cmpy(Node& a, Node& b) {
        return a.y<b.y; }</pre>
   void init(vector<pair<int,int>> &ip) {
     for (int i = 0; i < ssize(ip); i++)
  tie(tree[i].x, tree[i].y) = ip[i],</pre>
            tree[i].id = i:
     root = build(0, (int)ip.size()-1, 0);
   Node* build(int L, int R, int d) {
     if (L>R) return nullptr;
     int M = (L+R)/2;
     nth_element(tree+L,tree+M,tree+R+1,d
          %2?cmpy:cmpx);
     Node &o = tree[M]; o.f = d % 2;
o.x1 = o.x2 = o.x; o.y1 = o.y2 = o.y;
     o.L = build(L, M-1, d+1); o.R = build
     (M+1, R, d+1);
for (Node *s: {o.L, o.R}) if (s) {
       o.x1 = min(o.x1, s->x1); o.x2 = max
             (0.x2, s->x2);
       o.y1 = min(o.y1, s->y1); o.y2 = max
             (o.y2, s->y2);
     return tree+M:
   bool touch(int x, int y, lld d2, Node *
        r){
     lld d = (lld) \operatorname{sqrt}(d2) + 1;
     return x >= r->x1 - d & x <= r->x2 +
           d &&
                     y >= r -> y1 - d \&\& y <=
                           r->y^2+d;
   using P = pair<lld, int>;
   void dfs(int x, int y, P &mn, Node *r)
     if (!r || !touch(x, y, mn.first, r))
          return;
     mn = min(mn, P(dis2(r->x, r->y, x, y))
           , r->id));
     if (r->f == 1 ? y < r->y : x < r->x)
       dfs(x, y, mn, r\rightarrow L), dfs(x, y, mn,
            r->R);
     else
       dfs(x, y, mn, r\rightarrow R), dfs(x, y, mn,
             r->L);
   int query(int x, int y) {
     P mn(INF, -1); dfs(x, y, mn, root);
     return mn.second;
|} tree;
```

```
// from 8BQube, plz ensure p is strict convex hull
const llf INF = 1e18, qi = acos(-1) / 2 *
pair<llf, llf> solve(const vector<P> &p)
  llf mx = 0, mn = INF; int n = (int)p.
       size();
  for (int i = 0, u = 1, r = 1, l = 1; i
       < n; ++i) {
#define Z(v) (p[(v) % n] - p[i])
    P e = Z(i + 1);
    while (cross(e, Z(u + 1)) > cross(e,
        Z(u))) ++u;
    while (dot(e, Z(r + 1)) > dot(e, Z(r))
         )) ++r;
    if(!i) l = r + 1;
    while (dot(e, Z(l + 1)) < dot(e, Z(l))
        )) ++1;
    P D = p[r \% n] - p[l \% n];
    llf H = cross(e, Z(u)) / llf(norm(e))
    mn = min(mn, dot(e, D) * H);
    11f B = sqrt(norm(D)) * sqrt(norm(Z(u
         )));
    llf deg = (qi - acos(dot(D, Z(u)) / B
         )) / 2;
    mx = max(mx, B * sin(deg) * sin(deg))
  return {mn, mx};
} // test @ UVA 819
```

## 8.12 MinkowskiSum

```
// A, B are strict convex hull rotate to
      min by (X, Y)
 vector<P> Minkowski(vector<P> A, vector<P
     > B) {
   const int N = (int)A.size(), M = (int)B
        .size();
   vector<P> sa(N), sb(M), C(N + M + 1);
for (int i = 0; i < N; i++) sa[i] = A[(</pre>
        i+1)%N]-A[i];
   for (int i = 0; i < M; i++) sb[i] = B[(
        i+1)%M]-B[i];
   C[0] = A[0] + B[0];
   for (int i = 0, j = 0; i < N | | j < M;
        ) {
     P = (j>=M \mid i \in N \& cross(sa[i],
          sb[j])>=0))
       ? sa[i++] : sb[j++];
     C[i + j] = e;
   partial_sum(all(C), C.begin()); C.
        pop_back();
   return convex_hull(C); // just to
        remove colinear
} // be careful if min(|A|,|B|)<=2</pre>
```

#### 8.13 PointInHull

```
bool isAnti(P a, P b) {
  return cross(a, b) == 0 \& dot(a, b) <=
       0; }
bool PIH(const vector<P> &h, P z, bool
    strict = true) {
  int n = (int)h.size(), a = 1, b = n -
      1, r = !strict;
  if (n < 3) return r && isAnti(h[0] - z,</pre>
       h[n-1] - z);
  if (ori(h[0],h[a],h[b]) > 0) swap(a, b)
  if (ori(h[0],h[a],z) >= r || ori(h[0],h
       [b],z) <= -r)
    return false;
  while (abs(a - b) > 1) {
    int c = (a + b) / 2;
    (ori(h[0], h[c], z) > 0 ? b : a) = c;
  return ori(h[a], h[b], z) < r;</pre>
```

```
bool PIP(const vector<P> &p, P z, bool
    strict = true) {
    int cnt = 0, n = (int)p.size();
    for (int i = 0; i < n; i++) {
        P A = p[i], B = p[(i + 1) % n];
        if (isInter(Seg(A, B), z)) return !
            strict;
    auto zy = IM(z), Ay = IM(A), By = IM(
            B);
    cnt ^= ((zy<Ay) - (zy<By)) * ori(z, A
            , B) > 0;
    }
    return cnt;
}
```

## 8.15 PointInPolyFast

```
vector<int> PIPfast(vector<P> p, vector<P
    > q) {
  const int N = int(p.size()), Q = int(q.
       size());
  vector<pair<P, int>> evt; vector<Seg>
       edge;
  for (int i = 0; i < N; i++) {
    int a = i, b = (i + 1) \% N;
    P A = p[a], B = p[b];
assert (A < B || B < A); // std::
         operator<
    if (B < A) swap(A, B);
    evt.emplace_back(A, i); evt.
         emplace_back(B, ~i);
    edge.emplace_back(A, B);
  for (int i = 0; i < 0; i++)
    evt.emplace_back(q[i], i + N);
  sort(all(evt));
  auto vtx = p; sort(all(vtx));
auto eval = [](const Seg &a, lld x) ->
       11f {
    if (RE(a.dir) == 0) {
      assert (x == RE(a.st));
      return IM(a.st) + llf(IM(a.dir)) /
    llf t = (x - RE(a.st)) / llf(RE(a.dir))
         ));
    return IM(a.st) + IM(a.dir) * t;
  ild cur_x = 0;
  auto cmp = [&](const Seg &a, const Seg
       &b) -> bool {
    if (int s = sgn(eval(a, cur_x) - eval
         (b, cur_x)))
      return s == -1; // be careful: sgn<
           llf>, sgn<lld>
    int s = sgn(cross(b.dir, a.dir));
    if (cur_x != RE(a.st) && cur_x != RE(
         b.st)) s *= -1;
    return s == -1;
 };
namespace pbds = __gnu_pbds;
  pbds::tree<Seg, int, decltype(cmp),</pre>
    pbds::rb_tree_tag,
pbds::
         tree_order_statistics_node_update
         > st(cmp);
  auto answer = [\&](P ep) {
    if (binary_search(all(vtx), ep))
      return 1; // on vertex
    Seg H(ep, ep); // ??
    auto it = st.lower_bound(H);
    if (it != st.end() && isInter(it->
      first, ep))
return 1; // on edge
    if (it != st.begin() && isInter(prev())
         it)->first, ep))
      return 1; // on edge
    auto rk = st.order_of_key(H);
return rk % 2 == 0 ? 0 : 2; // 0:
         outside, 2: inside
  };
  vector<int> ans(Q);
  for (auto [ep, i] : evt) {
    cur_x = RE(ep);
if (i < 0) { // remove</pre>
      st.erase(edge[~i]);
    } else if (i < N) \{ // insert
```

## 8.11 MinMaxEnclosingRect 8.14 PointInPoly

## 8.16 PolyUnion

```
| llf polyUnion(const vector<vector<P>> &p)
   vector<tuple<P, P, int>> seg;
for (int i = 0; i < ssize(p); i++)</pre>
     for (int j = 0, m = int(p[i].size());
           j < m; j++)
       seg.emplace_back(p[i][j], p[i][(j +
             1) % m], i);
   llf ret = 0; // area of p[i] must be
        non-negative
   for (auto [A, B, i] : seg) {
     vector<pair<llf, int>> evt{\{0, 0\},
          {1, 0}};
     for (auto [C, D, j] : seg) {
  int sc = ori(A, B, C), sd = ori(A,
            B, D);
       if (sc != sd && i != j && min(sc,
            sd) < 0) {
         llf sa = cross(D-C, A-C), sb =
              cross(D-C, B-C);
         evt.emplace_back(sa / (sa - sb),
              sgn(sc - sd));
       } else if (!sc && !sd && j < i</pre>
           && sgn(dot(B - A, D - C)) > 0)
         evt.emplace_back(real((C - A) / (
              B - A)), 1);
         evt.emplace_back(real((D - A) / (
              B - A)), -1);
       }
     sort(evt.begin(), evt.end());
     llf sum = 0, last = 0; int cnt = 0;
     for (auto [q, c] : evt) {
       if (!cnt) sum += q - last;
cnt += c; last = q;
     ret += cross(A, B) * sum;
   return ret / 2;
i}
```

## 8.17 RotatingSweepLine

```
struct Event {
  Pd; int u, v;
  bool operator<(const Event &b) const {</pre>
    return sgn(cross(d, b.d)) > 0; }
P makePositive(P z) { return cmpxy(z, 0)
     ? -z : z; }
void rotatingSweepLine(const vector<P> &p
    ) {
  const int n = int(p.size());
  vector<Event> e; e.reserve(n * (n - 1)
       / 2);
  for (int i = 0; i < n; i++)
    for (int j = i + 1; j < n; j++)
      e.emplace_back(makePositive(p[i] -
           p[j]), i, j);
  sort(all(e));
  vector<int> ord(n), pos(n);
 iota(all(ord), 0);
sort(all(ord), [&p](int i, int j) {
  return cmpxy(p[i], p[j]); });
  for (int i = 0; i < n; i++) pos[ord[i]]</pre>
        = i;
  const auto makeReverse = [](auto &v) {
    sort(all(v)); v.erase(unique(all(v)),
          v.end());
    vector<pair<int,int>> segs;
    for (size_t i = 0, j = 0; i < v.size)
         (); i = j) {
      for (; j < v.size() && v[j] - v[i]</pre>
            <= j - i; j++);
```

```
segs.emplace_back(v[i], v[j - 1] +
         1 + 1);
  return segs;
for (size_t i = 0, j = 0; i < e.size();</pre>
  i = j) {
/* do here */
  vector<size_t> tmp;
  for (; j < e.size() && !(e[i] < e[j])</pre>
       ; j++)
    tmp.push_back(min(pos[e[j].u], pos[
         e[j].v]));
  for (auto [1, r] : makeReverse(tmp))
    reverse(ord.begin() + 1, ord.begin
         () + r);
    for (int t = 1; t < r; t++) pos[ord</pre>
         [t]] = t;
  }
}
```

### 8.18 SegIsIntersect

```
struct Seg { // closed segment
  P st, dir; // represent st + t*dir for
       0<=t<=1
  Seg(P s, P e) : st(s), dir(e - s) {}
  static bool valid(lld p, lld q) {
    // is there t s.t. 0 <= t <= 1 && qt
          == p ?
    if (q < 0) q = -q, p = -p;
    return sgn(0 - p) \le 0 \&\& sgn(p - q)
  vector<P> ends() const { return { st,
       st + dir }; }
template <typename T> bool isInter(T A, P
      p) {
  if (sgn(norm(A.dir)) == 0)
    return sgn(norm(p - A.st)) == 0; //
         BE CAREFUL
  return sgn(cross(p - A.st, A.dir)) == 0
        &&
    T::valid(dot(p - A.st, A.dir), norm(A
          .dir));
template <typename U, typename V>
bool isInter(U A, V B) {
  if (sgn(cross(A.dir, B.dir)) == 0) { //
        BE CAREFUL
    bool res = false;
    for (P p: A.ends()) res l= isInter(B,
    for (P p: B.ends()) res l= isInter(A,
    p);
return res;
  P D = B.st - A.st; lld C = cross(A.dir,
        B.dir);
  return U::valid(cross(D, B.dir), C) &&
    V::valid(cross(D, A.dir), C);
```

## 8.19 SegSegDist

```
// be careful of abs<complex<int>> (
    replace _abs below)
llf PointSegDist(P A, Seg B) {
  if (B.dir == P(0)) return _abs(A - B.st
      );
  if (sgn(dot(A - B.st, B.dir)) *
      sgn(dot(A - B.ed, B.dir)) <= 0)
    return abs(cross(A - B.st, B.dir)) /
         _abs(B.dir);
  return min(_abs(A - B.st), _abs(A - B.
      ed));
llf SegSegDist(const Seg &s1, const Seg &
    s2) {
  if (isInter(s1, s2)) return 0;
  return min({
      PointSegDist(s1.st, s2),
      PointSegDist(s1.ed, s2),
      PointSegDist(s2.st, s1),
```

```
PointSegDist(s2.ed, s1) });
|} // test @ QOJ2444 / PTZ19 Summer.D3
```

## 8.20 SimulateAnnealing

### 8.21 TangentPointToHull

#### 8.22 TriCenter

```
0 = ... // see min circle cover
G = (A + B + C) / 3;
H = G * 3 - 0 * 2; // orthogonal center
llf a = abs(B - C), b = abs(A - C), c =
    abs(A - B);
I = (a * A + b * B + c * C) / (a + b + c)
    ;
// FermatPoint: minimizes sum of distance
// if max. angle >= 120 deg then vertex
// otherwise, make eq. triangle AB'C, CA'
    B, BC'A
// line AA', BB', CC' intersects at P
```

#### 8.23 Voronoi

```
void build_voronoi_cells(auto &&p, auto
     &&res) {
   vector<vector<int>> adj(p.size());
   for (auto f: res) F3 {
     int a = f[i], b = f[R(i)];
     if (a >= p.size() || b >= p.size())
          continue:
     adj[a].emplace_back(b);
  // use `adj` and `p` and HPI to build
        cells
   for (size_t i = 0; i < p.size(); i++) {</pre>
     vector<Line> ls = frame; // the frame
     for (int j : adj[i]) {
       P m = p[i] + p[j], d = rot90(p[j] -
            p[i]);
       assert (norm(d) != 0);
       ls.emplace_back(m, m + d); //
           doubled coordinate
     } // HPI(ls)
| }
```

## 9 Miscellaneous

#### 9.1 Cactus 1

```
auto work = [&](const vector<int> cycle)
  // merge cycle info to u?
  int len = cycle.size(), u = cycle[0];
auto dfs = [&](auto dfs, int u, int p) {
  par[u] = p;
  vis[u] = 1;
  for (auto v : adj[u]) {
    if (v == p) { continue; }
    if (vis[v] == 0) {
      dfs(dfs, v, u);
      if (!cyc[v]) { // merge dp }
    } else if (vis[v] == 1) {
      for (int w = u; w != v; w = par[w])
        cyc[w] = 1;
      }
    } else {
      vector<int> cycle = {u};
      for (int w = v; w != u; w = par[w])
            { cycle.push_back(w); }
      work(cycle);
  vis[u] = 2;
|};
```

#### 9.2 Cactus 2

```
// a component contains no articulation
     point, so P2 is a component
// but not a vertex biconnected component
      by definition
// resulting bct is rooted
struct BlockCutTree {
  int n, square = 0, cur = 0;
  vector<int> low, dfn, stk;
  vector<vector<int>> adj, bct;
  BlockCutTree(int n) : n(n), low(n), dfn
  (n, -1), adj(n), bct(n) {}
void build() { dfs(0); }
void addEdge(int u, int v) { adj[u].
        push_back(v), adj[v].push_back(u);
  void dfs(int u) {
    low[u] = dfn[u] = cur++;
     stk.push_back(u);
     for (auto v : adj[u]) {
       if (dfn[v] == -1) {
         dfs(v);
         low[u] = min(low[u], low[v]);
         if (low[v] == dfn[u]) {
           bct.emplace_back();
           int x;
           do {
             x = stk.back();
             stk.pop_back();
             bct.back().push_back(x);
           } while (x != v);
           bct[u].push_back(n + square);
           sauare++;
         }
       } else {
         low[u] = min(low[u], dfn[v]);
       }
    }
  }
j};
```

## 9.3 Dancing Links

```
|#include <bits/stdc++.h>
|using namespace std;
|// tioj 1333
|#define TRAV(i, link, start) for (int i = link[start]; i != start; i = link[i ])
|const int NN = 40000, RR = 200;
|template<br/>
| template<br/>
| 1s, RR: num of rows
```

```
struct DLX {
 int lt[NN], rg[NN], up[NN], dn[NN], rw[
       NN], cl[NN], bt[NN], s[NN], head,
       sz, ans;
  int rows, columns;
  bool vis[NN];
 bitset<RR> sol, cur; // not sure
 void remove(int c) {
    if (E) lt[rg[c]] = lt[c], rg[lt[c]] =
          rg[c];
    TRAV(i, dn, c) {
      if (E) {
        TRAV(j, rg, i)
up[dn[j]] = up[j], dn[up[j]] =
                dn[j], --s[cl[j]];
        lt[rg[i]] = lt[i], rg[lt[i]] = rg
             Γi];
    }
 void restore(int c) {
   TRAV(i, up, c) {
   if (E) {
        TRAV(j, lt, i)
          ++s[cl[j]], up[dn[j]] = j, dn[
               up[j]] = j;
      } else {
        lt[rg[i]] = rg[lt[i]] = i;
      }
    if (E) lt[rg[c]] = c, rg[lt[c]] = c;
 void init(int c) {
  rows = 0, columns = c;
  for (int i = 0; i < c; ++i) {</pre>
      up[i] = dn[i] = bt[i] = i;
      lt[i] = i == 0 ? c : i - 1;
      rg[i] = i == c - 1 ? c : i + 1;
      s[i] = 0;
    rg[c] = 0, lt[c] = c - 1;
    up[c] = dn[c] = -1;
    head = c, sz = c + 1;
 void insert(const vector<int> &col) {
    if (col.empty()) return;
    int f = sz;
    for (int i = 0; i < (int)col.size();</pre>
         ++i) {
      int c = col[i], v = sz++;
      dn[bt[c]] = v;
      up[v] = bt[c], bt[c] = v;
      rg[v] = (i + 1 == (int)col.size() ?
            f: v + 1);
      rw[v] = rows, cl[v] = c;
      ++s[c];
      if (i > 0) lt[v] = v - 1;
    ++rows, lt[f] = sz - 1;
  int h() {
    int ret = 0;
    fill_n(vis, sz, false);
TRAV(x, rg, head) {
      if (vis[x]) continue;
      vis[x] = true, ++ret;
TRAV(i, dn, x) TRAV(j, rg, i) vis[
           cl[j]] = true;
    return ret;
 void dfs(int dep) {
    if (dep + (E ? 0 : h()) >= ans)
    if (rg[head] == head) return sol =
         cur, ans = dep, void();
    if (dn[rg[head]] == rg[head]) return;
    int w = rg[head];
    TRAV(x, rg, head) if (s[x] < s[w]) w
         = x;
    if (E) remove(w);
    TRAV(i, dn, w) {
      if (!E) remove(i);
      TRAV(j, rg, i) remove(E ? cl[j] : j
```

```
cur.set(rw[i]), dfs(dep + 1), cur.
            reset(rw[i]);
       TRAV(j, lt, i) restore(E ? cl[j] :
       if (!E) restore(i);
    if (E) restore(w);
  }
  int solve() {
    for (int i = 0; i < columns; ++i)</pre>
      dn[bt[i]] = i, up[i] = bt[i];
    ans = 1e9, sol.reset(), dfs(0);
    return ans;
  }
};
int main() {
    int n, m; cin >> n >> m;
DLX<true> solver;
    solver.init(m);
    for (int i = 0; i < n; i++){
         vector<int> add;
         for (int j = 0; j < m; j++){
             int x; cin >> x;
if (x == 1) {
                  add.push_back(j);
         solver.insert(add);
    cout << solver.solve() << '\n';</pre>
    return 0;
```

## 9.4 Offline Dynamic MST

```
int cnt[maxn], cost[maxn], st[maxn], ed[
     maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed
      , qr[i].second = weight after
     operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v
     contains edges i such that cnt[i] ==
void contract(int 1, int r, vector<int> v
      vector<int> &x, vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int
      if (cost[i] == cost[j]) return i <</pre>
           j;
      return cost[i] < cost[j];</pre>
      });
  djs.save();
  for (int i = l; i <= r; ++i) djs.merge(</pre>
       st[qr[i].first], ed[qr[i].first]);
  for (int i = 0; i < (int)v.size(); ++i)</pre>
    if (djs.find(st[v[i]]) != djs.find(ed
         [v[i]])) {
      x.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
    }
  }
  djs.undo();
  dis.save();
  for (int i = 0; i < (int)x.size(); ++i)
        djs.merge(st[x[i]], ed[x[i]]);
  for (int i = 0; i < (int)v.size(); ++i)</pre>
    if (djs.find(st[v[i]]) != djs.find(ed
         [v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
    }
  djs.undo();
}
void solve(int l, int r, vector<int> v,
     long long c) {
  if (l == r) {
    cost[qr[l].first] = qr[l].second;
    if (st[qr[l].first] == ed[qr[l].first
```

7) -

printf("%lld\n", c);

```
return:
     int minv = qr[l].second;
     for (int i = 0; i < (int)v.size(); ++</pre>
          i) minv = min(minv, cost[v[i]]);
     printf("%lld\n", c + minv);
     return:
  int m = (l + r) >> 1;
  vector < int > lv = v, rv = v;
   vector<int> x, y;
   for (int i = m + 1; i \le r; ++i) {
     cnt[qr[i].first]--;
     if (cnt[qr[i].first] == 0) lv.
          push_back(qr[i].first);
  contract(l, m, lv, x, y);
long long lc = c, rc = c;
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i)</pre>
     lc += cost[x[i]];
     djs.merge(st[x[i]], ed[x[i]]);
  solve(l, m, y, lc);
  djs.undo();
   x.clear(), y.clear();
   for (int i = m + 1; i <= r; ++i) cnt[qr</pre>
        [i].first]++;
  for (int i = 1; i <= m; ++i) {</pre>
     cnt[qr[i].first]--;
     if (cnt[qr[i].first] == 0) rv.
          push_back(qr[i].first);
  contract(m + 1, r, rv, x, y);
  djs.save();
   for (int i = 0; i < (int)x.size(); ++i)</pre>
         {
     rc += cost[x[i]];
     djs.merge(st[x[i]], ed[x[i]]);
  solve(m + 1, r, y, rc);
  djs.undo();
  for (int i = 1; i <= m; ++i) cnt[qr[i].</pre>
        first]++;
į }
```

#### 9.5 Matroid Intersection

```
    x → y if S - {x} ∪ {y} ∈ I<sub>1</sub> with cost({y}).
    source → y if S ∪ {y} ∈ I<sub>1</sub> with cost({y}).
    y → x if S - {x} ∪ {y} ∈ I<sub>2</sub> with -cost({y}).
    y → sink if S ∪ {y} ∈ I<sub>2</sub> with -cost({y}).
```

Augmenting path is shortest path from source to sink.

#### 9.6 Euler Tour

```
vector<int> euler, vis(V);
auto dfs = [&](auto dfs, int u) -> void {
  while (!adj[u].empty()) {
    while (!adj[u].empty()) && del[adj[u].
        back()[1]]) {
        adj[u].pop_back();
    }
    if (!adj[u].empty()) {
        auto [v, i] = adj[u].back();
        del[i] = true;
        dfs(dfs, v);
    }
    euler.push_back(u);
};
dfs(dfs, 0);
reverse(euler.begin(), euler.end());
```

#### 9.7 SegTree Beats

```
| struct SegmentTree {
    int n;
    struct node {
        i64 mx1, mx2, mxc;
        i64 mn1, mn2, mnc;
        i64 add;
        i64 sum;
        node(i64 v = 0) {
            mx1 = mn1 = sum = v;
    }
```

```
mxc = mnc = 1;
    add = 0;
    mx2 = -9e18, mn2 = 9e18;
 }
vector<node> t;
// build
void push(int id, int l, int r) {
  auto& c = t[id];
  int m = l + r \gg 1;
  if (c.add != 0) {
    apply_add(id << 1, 1, m, c.add);</pre>
    apply_add(id \ll 1 | 1, m + 1, r, c.
         add);
    c.add = 0:
  apply_min(id << 1, 1, m, c.mn1);</pre>
  apply_min(id \ll 1 | 1, m + 1, r, c.
       mn1):
  apply_max(id \ll 1, l, m, c.mx1);
  apply_max(id \ll 1 | 1, m + 1, r, c.
void apply_add(int id, int l, int r,
    i64 v) {
  if (v == 0) {
    return;
  auto& c = t[id];
  c.add += v;
  c.sum += v * (r - l + 1);
  c.mx1 += v;
  c.mn1 += v;
  if (c.mx2 != -9e18) {
    c.mx2 += v;
  if (c.mn2 != 9e18) {
    c.mn2 += v;
  }
void apply_min(int id, int l, int r,
     i64 v) {
  auto& c = t[id];
  if (v <= c.mn1) {</pre>
    return;
  c.sum -= c.mn1 * c.mnc;
  c.mn1 = v;
  c.sum += c.mn1 * c.mnc;
  if (l == r || v >= c.mx1) {
    c.mx1 = v;
  } else if (v > c.mx2) {
    c.mx2 = v;
 }
void apply_max(int id, int l, int r,
     i64 v) {
  auto& c = t[id];
  if (v >= c.mx1) {
    return;
  c.sum -= c.mx1 * c.mxc;
  c.mx1 = v;
  c.sum += c.mx1 * c.mxc;
  if (l == r \mid \mid v \Leftarrow c.mn1) {
    c.mn1 = v;
  } else if (v < c.mn2) {
    c.mn2 = v;
 }
void pull(int id) {
  auto &c = t[id], &lc = t[id << 1], &</pre>
       rc = t[id << 1 | 1];
  c.sum = lc.sum + rc.sum;
  if (lc.mn1 == rc.mn1) {
    c.mn1 = lc.mn1;
    c.mn2 = min(lc.mn2, rc.mn2);
    c.mnc = lc.mnc + rc.mnc;
 } else if (lc.mn1 < rc.mn1) {
   c.mn1 = lc.mn1;</pre>
    c.mn2 = min(lc.mn2, rc.mn1);
    c.mnc = lc.mnc;
  } else {
    c.mn1 = rc.mn1;
    c.mn2 = min(lc.mn1, rc.mn2);
    c.mnc = rc.mnc;
  if (lc.mx1 == rc.mx1) {
    c.mx1 = lc.mx1;
```

```
c.mx2 = max(lc.mx2, rc.mx2);
c.mxc = lc.mxc + rc.mxc;
  } else if (lc.mx1 > rc.mx1) {
   c.mx1 = lc.mx1;
    c.mx2 = max(lc.mx2, rc.mx1);
    c.mxc = lc.mxc;
  } else {
    c.mx1 = rc.mx1;
    c.mx2 = max(lc.mx1, rc.mx2);
    c.mxc = rc.mxc:
  }
void range_chmin(int id, int l, int r,
  int ql, int qr, i64 v) {
if (r < ql || l > qr || v >= t[id].
       mx1) {
    return:
  if (ql <= l && r <= qr && v > t[id].
       mx2) {
    apply_max(id, l, r, v);
  push(id, l, r);
  int m = l + r >> 1;
  range_chmin(id << 1, l, m, ql, qr, v)</pre>
  range_chmin(id \ll 1 | 1, m + 1, r, ql
       , qr, v);
  pull(id);
void range_chmin(int ql, int qr, i64 v)
  range_chmin(1, 0, n - 1, ql, qr, v);
}
void range_chmax(int id, int l, int r,
     int ql, int qr, i64 v) {
  if (r < ql \mid | l > qr \mid | v \leftarrow t[id].
       mn1) {
    return;
  if (ql \le l \& r \le qr \& v < t[id].
       mn2) {
    apply_min(id, l, r, v);
    return;
  push(id, l, r);
  int m = 1 + r >> 1;
  range_chmax(id \ll 1, l, m, ql, qr, v)
  range\_chmax(id << 1 | 1, m + 1, r, ql)
         qr, v);
  pull(id);
void range_chmax(int ql, int qr, i64 v)
  range_chmax(1, 0, n - 1, ql, qr, v);
void range_add(int id, int l, int r,
  int ql, int qr, i64 v) {
if (r < ql || l > qr) {
    return:
  if (ql <= l && r <= qr) {</pre>
    apply_add(id, l, r, v);
    return;
  push(id, l, r);
int m = l + r >> 1;
  range_add(id << 1, l, m, ql, qr, v);
range_add(id << 1 | 1, m + 1, r, ql,
       qr, v);
  pull(id);
void range_add(int ql, int qr, i64 v) {
  range_add(1, 0, n - 1, ql, qr, v);
i64 range_sum(int id, int l, int r, int
      ql, int qr) {
  if (r < ql || l > qr) {
   return 0;
  if (ql \ll l \& r \ll qr) {
    return t[id].sum;
  push(id, 1, r);
  int m = l + r \gg 1;
```

#### 9.8 Decimal

## 9.9 AdaSimpson

```
template<typename Func, typename d =
     double>
struct Simpson {
  using pdd = pair<d, d>;
  Func f:
  pdd mix(pdd l, pdd r, optional<d> fm =
       {}) {
    d h = (r.X - 1.X) / 2, v = fm.
          value_or(f(1.X + h));
    return \{v, h / 3 * (1.Y + 4 * v + r.Y)\}
          )};
 d eval(pdd l, pdd r, d fm, d eps) {
  pdd m((l.X + r.X) / 2, fm);
    d s = mix(1, r, fm).second;
    auto [flm, sl] = mix(l, m);
auto [fmr, sr] = mix(m, r);
d delta = sl + sr - s;
if (abs(delta) <= 15 * eps) return sl
+ sr + delta / 15;
    return eval(1, m, flm, eps / 2) +
       eval(m, r, fmr, eps / 2);
 d = val2(d l, d r, d eps, int k = 997) {
    d h = (r - 1) / k, s = 0;
    for (int i = 0; i < k; ++i, l += h)
s += eval(l, l + h, eps / k);
    return s;
 }
template<typename Func>
Simpson<Func> make_simpson(Func f) {
     return {f}; }
```

## 9.10 SB Tree

```
struct Q {
  11 p, q;
  Q go(Q b, ll d) { return \{p + b.p*d, q\}
        + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such
     that
// pred(p/q) is true, and 0 \le p,q \le N
Q frac_bs(ll N) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
 assert(pred(hi));
bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step/=2)
            step*=2);)
       if (Q mid = hi.go(lo, len + step);
   mid.p > N || mid.q > N || dir ^
                  pred(mid))
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
```

```
|  }
| return dir ? hi : lo;
|}
```

#### 9.11 Bitset LCS

```
cin >> n >> m;
for (int i = 1, x; i <= n; ++i)
  cin >> x, p[x].set(i);
for (int i = 1, x; i <= m; i++) {
  cin >> x, (g = f) |= p[x];
  f.shiftLeftByOne(), f.set(0);
  ((f = g - f) ^= g) &= g;
}
cout << f.count() << '\n';</pre>
```

#### 9.12 Hilbert Curve

#### 9.13 Mo on Tree

```
void MoAlgoOnTree() {
 Dfs(0, -1);
 vector<int> euler(tk);
  for (int i = 0; i < n; ++i) {
    euler[tin[i]] = i;
   euler[tout[i]] = i;
 vector<int> l(q), r(q), qr(q), sp(q)
      -1);
  for (int i = 0; i < q; ++i) {
    if (tin[u[i]] > tin[v[i]]) swap(u[i],
         v[i]);
   int z = GetLCA(u[i], v[i]);
    sp[i] = z[i];
    if (z == u) l[i] = tin[u[i]], r[i] =
        tin[v[i]];
    else l[i] = tout[u[i]], r[i] = tin[v[
        i]];
    qr[i] = i;
 sort(qr.begin(), qr.end(), [&](int i,
      int j) {
      if (l[i] / kB == l[j] / kB) return
          r[i] < r[j];
      return l[i] / kB < l[j] / kB;</pre>
     });
 vector<bool> used(n);
  // Add(v): add/remove v to/from the
      path based on used[v]
  for (int i = 0, tl = 0, tr = -1; i < q;
        ++i) {
    while (tl < l[qr[i]]) Add(euler[tl</pre>
        ++]);
    while (tl > l[qr[i]]) Add(euler[--tl
        ]);
   while (tr > r[qr[i]]) Add(euler[tr
        --1):
    while (tr < r[qr[i]]) Add(euler[++tr</pre>
    // add/remove LCA(u, v) if necessary
```

## 9.14 N Queens

```
void solve(vector<int> &ret, int n) { //
no sol when n=2,3
  if (n % 6 == 2) {
    for (int i = 2; i <= n; i += 2) ret.
        pb(i);
    ret.pb(3); ret.pb(1);</pre>
```

```
for (int i = 7; i \le n; i += 2) ret.
          pb(i);
     ret.pb(5);
  } else if (n % 6 == 3) {
     for (int i = 4; i <= n; i += 2) ret.
          pb(i);
     ret.pb(2);
     for (int i = 5; i \le n; i += 2) ret.
          pb(i);
     ret.pb(1); ret.pb(3);
  } else {
     for (int i = 2; i \leftarrow n; i \leftarrow 2) ret.
          pb(i);
     for (int i = 1; i \le n; i += 2) ret.
          pb(i);
  }
į }
```

#### 9.15 Rollback Mo

```
for (int l = 0, r = -1; auto [ql, qr, i]
      : qs) {
  if (ql / B == qr / B) {
    for (int j = ql; j <= qr; j++) {</pre>
      cntSmall[a[j]]++;
      ans[i] = max(ans[i], 1LL * b[a[j]]
             * cntSmall[a[j]]);
    for (int j = ql; j <= qr; j++) {</pre>
      cntSmall[a[j]]--;
    }
    continue;
  if (int block = ql / B; block != lst) {
    int x = min((block + 1) * B, n);
    while (r + 1 < x) { add(++r); }
while (r >= x) { del(r--); }
    while (l < x) { del(l++); }</pre>
    mx = 0;
    lst = block;
  while (r < qr) \{ add(++r); \}
  i64 \text{ tmpMx} = mx;
  int tmpL = 1:
  while (l > ql) { add(--l); }
  ans[i] = mx;
  mx = tmpMx;
  while (l < tmpL) { del(l++); }</pre>
```

## 9.16 Subset Sum

```
template<size_t S> // sum(a) < S
bitset<S> SubsetSum(const int *a, int n)
    {
      vector<int> c(S);
      bitset<S> dp; dp[0] = 1;
      for (int i = 0; i < n; ++i) ++c[a[i ]];
      for (size_t i = 1; i < S; ++i) {
            while (c[i] > 2) c[i] -= 2, ++c[i * 2];
            while (c[i]--) dp |= dp << i;
      }
      return dp;
}</pre>
```