

This module contains *TLAPS* checked proofs of the safety properties asserted in module *ParReach*—namely, the invariance of *Inv* and that the parallel algorithm implements the safety part of *Misra*’s algorithm under the refinement mapping defined there.

EXTENDS *ParReach*, *Integers*, *TLAPS*

LEMMA *TypeInvariant* \triangleq *Spec* $\Rightarrow \Box$ *Inv*

$\langle 1 \rangle 1$. *Init* \Rightarrow *Inv*

BY *RootAssump* DEF *Init*, *Inv*, *ProcSet*

$\langle 1 \rangle 2$. *Inv* $\wedge [Next]_{vars} \Rightarrow Inv'$

BY *SuccAssump* DEF *Inv*, *Next*, *vars*, *ProcSet*, *p*, *a*, *b*, *c*

$\langle 1 \rangle 3$. QED

BY $\langle 1 \rangle 1$, $\langle 1 \rangle 2$, *PTL* DEF *Spec*

THEOREM *Spec* $\Rightarrow R!Init \wedge \Box[R!Next]_R!vars$

$\langle 1 \rangle 1$. *Init* $\Rightarrow R!Init$

BY *ProcsAssump* DEF *Init*, *R!Init*, *pcBar*, *vrootBar*, *ProcSet*

$\langle 1 \rangle 2$. *Inv* $\wedge [Next]_{vars} \Rightarrow [R!Next]_R!vars$

$\langle 2 \rangle$ SUFFICES ASSUME *Inv*,

PROVE $\frac{[Next]_{vars}}{[R!Next]_R!vars}$

OBVIOUS

$\langle 2 \rangle$ USE DEF *Inv*, *Next*, *vars*, *R!Next*, *R!vars*, *vrootBar*, *pcBar*

$\langle 2 \rangle 1$. ASSUME NEW *self* \in *Procs*,

a(*self*)

PROVE $[R!Next]_R!vars$

$\langle 3 \rangle 1$. ASSUME *vroot* $\neq \{\}$

PROVE UNCHANGED *R!vars*

BY $\langle 2 \rangle 1$, $\langle 3 \rangle 1$ DEF *a*

$\langle 3 \rangle 2$. ASSUME *vroot* $= \{\}$

PROVE $[R!Next]_R!vars$

$\langle 4 \rangle 1$. ASSUME *vrootBar* $= \{\}$

PROVE $[R!Next]_R!vars$

BY $\langle 2 \rangle 1$, $\langle 3 \rangle 2$, $\langle 4 \rangle 1$ DEF *a*, *R!a*

$\langle 4 \rangle 2$. ASSUME *vrootBar* $\neq \{\}$

PROVE UNCHANGED *R!vars*

$\langle 5 \rangle 1$. $\exists q \in Procs \setminus \{self\} : pc[q] \neq \text{“Done”}$

BY $\langle 4 \rangle 2$, $\langle 3 \rangle 2$, $\langle 2 \rangle 1$ DEF *a*

$\langle 5 \rangle 2$. *pcBar'* $\neq \text{“Done”}$

BY $\langle 5 \rangle 1$, $\langle 3 \rangle 2$, $\langle 2 \rangle 1$ DEF *a*

$\langle 5 \rangle$. QED

BY $\langle 5 \rangle 2$, $\langle 3 \rangle 2$, $\langle 2 \rangle 1$ DEF *a*

$\langle 4 \rangle 3$. QED

BY $\langle 4 \rangle 1$, $\langle 4 \rangle 2$

$\langle 3 \rangle 3$. QED

BY $\langle 3 \rangle 1$, $\langle 3 \rangle 2$ DEF *R!Next*

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    <2>2. ASSUME NEW self ∈ Procs,
        b(self)
        PROVE [R!Next]R!vars
        BY <2>2 DEF b, R!a
    <2>3. ASSUME NEW self ∈ Procs,
        c(self)
        PROVE [R!Next]R!vars
        BY <2>3 DEF c
    <2>4.CASE UNCHANGED vars
        BY <2>4
    <2>5. QED
        BY <2>1, <2>2, <2>3, <2>4 DEF Next, p
    <1>3. QED
        BY <1>1, <1>2, TypeInvariant, PTL DEF Spec

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\ * Modification History
\ * Last modified Sun Apr 14 16:55:36 PDT 2019 by lamport
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