

# Homework 3

## 10-405/605: Machine Learning with Large Datasets

Due Wednesday, March 15th at 11:59 PM Eastern Time

Submit your solutions via Gradescope, **with your solution to each subproblem on a separate page**, i.e., following the template below.

**IMPORTANT:** Be sure to highlight where your solutions are for each question when submitting to Gradescope otherwise you will be marked 0 and will need to submit regrade request for every solution un-highlighted in order for fix it!

Note that Homework 3 consists of two parts: this written assignment, and a programming assignment. Remember to fill out the collaboration section found at the end of this homework as per the course policy.

**10-405 Students** are required to complete the following:

1. Written Section 1.1 (a) and (b) *[15 points]*
2. Written Section 1.2 (a) and (b) *[16 points]*
3. Written Section 1.3 (a), (b), (c), (d), and (e) *[15 points]*
4. All of Section 2 *[40 points]*

Your homework score will be calculated as a percentage over the maximum points you can earn, which is 86.

**10-605 Students** are required to complete **all** sections of the homework. Your homework score will be calculated as a percentage over the maximum points you can earn, which is 100.

**Programming:** The programming in this homework is **NOT** autograded however you are required to upload your completed notebooks to Gradescope, otherwise you will not receive credit for the programming sections.

# 1 Written Section [60 Points]

## 1.1 Count-Min Sketch Basics (15 pts)

- (a) Let's use count-min sketch to estimate the counts of a stream of 5 items:  $x_1, x_2, x_3, x_4, x_5$ , using 3 hash functions:  $h_1, h_2, h_3$ , each with 3 buckets. The hash functions have the following collision matrix:

collision matrix			
$h_1$	$x_1, x_3$	$x_4, x_5$	$x_2$
$h_2$	$x_2, x_5$	$x_3$	$x_1, x_4$
$h_3$	$x_1$	$x_2, x_3$	$x_4, x_5$

Suppose you see the following stream of data:  $\{x_1, x_2, x_5, x_3, x_3, x_4, x_4, x_4\}$ .

- [2 points]* What is the value of  $\hat{c}_1$ , the approximate count of item  $x_1$ , based on this sketch?
  - [2 points]* What is the value of  $\hat{c}_4$ , the approximate count of item  $x_4$ , based on this sketch?
  - [5 points]* Let  $c_s$  be the true count of item  $x_s$ . What is the average absolute error of approximation for all items, i.e.  $\frac{1}{5} \sum_{s=1}^5 |\hat{c}_s - c_s|$ ?
- (b) Consider the approximation guarantee for count-min sketch presented in lecture: given  $r > \log_2(\frac{1}{\delta})$  hash functions, each with  $m > \frac{2}{\epsilon}$  buckets, then

$$P(\hat{c}_s \geq c_s + \epsilon \|c\|_1) \leq \delta$$

- [3 points]* How would the number of required buckets change if our tolerance on the discrepancy  $\epsilon$  is halved?
- [3 points]* How would the number of hash functions change if we allowed the probability  $\delta$  to be twice as large?

## 1.2 An Unbiased Sketch (30 pts)

The count-min sketch presented in lecture is biased, as the estimated count  $\hat{c}_s$  for item  $s \in \{1, \dots, k\}$  is always higher than (or equal to) the true count  $c_s$ . In this problem, we will explore an *unbiased* sketch.

Let  $g$  be a hash function chosen from a family  $G$  of independent hashes, such that  $g$  maps each item  $s$  to either  $+1$  or  $-1$  with equal probability:

$$P(g(s) = +1) = P(g(s) = -1) = 1/2.$$

Let's start with a simple version of the sketch, which uses one additional hash function,  $h$ , which maps each item  $s$  to one of  $m$  buckets. Assume that  $g$  and  $h$  are independent. When we observe an item  $s$  in the sequence, we simply update the sketch counter,  $C$ , as:

$$C[h(s)] = C[h(s)] + g(s).$$

Then, if we would like to predict the count for item  $s$ , we return:

$$\hat{c}_s = C[h(s)] g(s).$$

Given this sketch, please answer the following questions:

- a) *[8 points]* First, show that this simple sketch is unbiased, i.e.,  $\mathbb{E}[\hat{c}_s] = c_s$ . (*Hint: Find an expression for  $\hat{c}_s$  in terms of  $g(s)$  and  $c_s$ , and use linearity of expectation.*)

- b) *[8 points]* Let's understand how much the estimates vary. Prove that:  $Var(\hat{c}_s) = \frac{1}{m} \sum_{j \in \{1, \dots, k\}, j \neq s} c_j^2$ , where  $m$  is the number of buckets for the hash function  $h$ .

**\*\*10-605 Students Only\*\***

- c) *[6 points]* We will now bound the probability of getting a bad estimate  $\hat{c}_s$ , i.e., one that differs substantially from the true count  $c_s$ . Show that:

$$P(|\hat{c}_s - c_s| \geq \epsilon \|\mathbf{c}_{-s}\|_2) \leq \frac{1}{\epsilon^2 m},$$

where  $\mathbf{c}_{-s}$  is a vector of length  $(k - 1)$  containing the counts for all items except  $s$ . (*Hint: Consider Chebyshev's inequality.*)

**\*\*10-605 Students Only\*\***

- d) Similar to count-min sketch, we can improve the results from (c) by using  $r$  hash functions for mapping the items to buckets  $(h_1, \dots, h_r)$ , and  $r$  hash functions for assigning signs  $(g_1, \dots, g_r)$ . In particular, setting  $m > \frac{3}{\epsilon^2}$  and  $r > (\log(1/\delta))$ , we can guarantee:

$$P(|\hat{c}_s - c_s| \geq \epsilon \|\mathbf{c}_{-s}\|_2) \leq \delta.$$

Note that while this result appears similar to the one we had for count-min sketch, it requires more buckets,  $m$ . Recall that for count-min sketch with  $m > \frac{2}{\epsilon}$  and  $r > (\log(1/\delta))$ , we had the bound:

$$P(|\hat{c}_s - c_s| \geq \epsilon \|\mathbf{c}\|_1) \leq \delta.$$

With these results in mind, answer the questions below.

- i. *[4 points]* Explain a scenario where count-min sketch would be preferable to the unbiased sketch we derived (using the above bounds).
  
  
  
  
  
  
  
  
  
  
- ii. *[4 points]* Explain a scenario where the unbiased sketch we derived would be preferable to count-min sketch (using the above bounds).

### 1.3 Locality-Sensitive Hashing (15 pts)

In lecture, we saw that LSH can be used for applications such as nearest neighbors classification and clustering. We mentioned that if two points  $\mathbf{x}$  and  $\mathbf{y}$  have small *cosine distance*, then the resulting bit vectors formed by using LSH,  $\mathbf{x}'$  and  $\mathbf{y}'$ , will have small *Hamming distance*. In this problem, we will explore this result in more detail.

First, let's recall the definitions of cosine similarity, cosine distance, and Hamming distance.

*Cosine similarity* measures the cosine of the angle between two vectors  $\mathbf{x}, \mathbf{y} \in \mathbb{R}^k$ , and can be calculated as:

$$s_{\cos}(\mathbf{x}, \mathbf{y}) = \cos(\theta(\mathbf{x}, \mathbf{y})) = \frac{\mathbf{x} \cdot \mathbf{y}}{\|\mathbf{x}\|_2 \|\mathbf{y}\|_2}$$

*Cosine distance* is measured as  $d_{\cos}(\mathbf{x}, \mathbf{y}) = 1 - s_{\cos}(\mathbf{x}, \mathbf{y})$ .

*Hamming distance* between two *binary* (0-1) vectors  $\mathbf{x}', \mathbf{y}'$  of length  $k$  measures the number of differing bits, and is calculated as:

$$d_h(\mathbf{x}', \mathbf{y}') = \sum_{i=1}^k |\mathbf{x}'_i - \mathbf{y}'_i|$$

(a) [1 points] Define  $\mathbf{a} = [0, 0, 1, 1, 0]$ ,  $\mathbf{b} = [0, 1, 0, 0, 1]$ . What's the *cosine similarity* between  $\mathbf{a}$  and  $\mathbf{b}$ ?

(b) [1 points] Define  $\mathbf{a} = [0, 0, 1, 1, 0]$ ,  $\mathbf{b} = [0, 1, 0, 0, 1]$ . What's the *Hamming distance* between  $\mathbf{a}$  and  $\mathbf{b}$ ?

(c) *[3 points]* For a random hyperplane,  $\mathbf{z}$ , define the following hash function:

$$h_z(\mathbf{x}) = \begin{cases} 1 & \text{if } \mathbf{z} \cdot \mathbf{x} \geq 0 \\ 0 & \text{if } \mathbf{z} \cdot \mathbf{x} < 0 \end{cases}$$

For two vectors  $\mathbf{x}, \mathbf{y} \in \mathbb{R}^k$ , prove that:

$$P[h_z(\mathbf{x}) = h_z(\mathbf{y})] = 1 - \frac{\theta(\mathbf{x}, \mathbf{y})}{\pi},$$

where  $\theta(\mathbf{x}, \mathbf{y})$  is the angle between vectors  $\mathbf{x}, \mathbf{y}$  measured in radians.

To prove this, use the fact that  $P[\mathbf{z} \cdot \mathbf{x} \geq 0, \mathbf{z} \cdot \mathbf{y} < 0] = \frac{\theta(\mathbf{x}, \mathbf{y})}{2\pi}$ , i.e., the probability that a random hyperplane separates two vectors is proportional to the angle between them.



- (d) *[6 points]* In LSH, we estimate  $P[h_z(\mathbf{x}) = h_z(\mathbf{y})]$  by selecting  $b$  random hyperplanes, and calculating bit vectors  $\mathbf{x}'_i = h_i(\mathbf{x})$  and  $\mathbf{y}'_i = h_i(\mathbf{y})$  for  $i = 1, \dots, b$ . Suppose that after running this procedure, we compute the Hamming distance and find:  $d_h(\mathbf{x}', \mathbf{y}') = h$ . Using this estimate and part (c), show that:

$$d_{cos}(\mathbf{x}, \mathbf{y}) \approx 1 - \cos\left(\frac{h}{b}\pi\right)$$

Note that this equation shows that if  $h = b$ , i.e., the bit vectors have a large hamming distance, the cosine distance will also be large ( $d_{cos}(\mathbf{x}, \mathbf{y}) = 2$ ). Similarly, for  $h = 0$ , we will have  $d_{cos}(\mathbf{x}, \mathbf{y}) = 0$ .  
*(Hint: The proof should only be a few lines.)*

(e) Let's go through a simple example using LSH. Suppose we have two vectors,  $\mathbf{x} = [3, 4, 5, 6]$  and  $\mathbf{y} = [4, 3, 2, 1]$ .

i. *[1 points]* Use one random vector:  $\mathbf{z}_1 = [0.16, -0.26, -1.12, 0.01]$  to perform LSH. Compute the error of the resulting estimated cosine distance, i.e.,  $|1 - \cos(\frac{b}{b} \pi) - d_{\cos}(\mathbf{x}, \mathbf{y})|$ .

ii. *[1 points]* Now let's try with more vectors. Use four random vectors:  $\mathbf{z}_1 = [0.16, -0.26, -1.12, 0.01]$ ,  $\mathbf{z}_2 = [1.08, -1.14, 0.8, -0.8]$ ,  $\mathbf{z}_3 = [1.94, 0.08, 1.83, 0.17]$ ,  $\mathbf{z}_4 = [-0.88, -1.01, 1.36, -0.07]$  to perform LSH. Compute the error of the estimated cosine distance.

iii. *[2 points]* In practice, it can be sufficient to consider random vectors only consisting of  $\{-1, +1\}$  when performing LSH. Use the signs of the previous random vectors, i.e.,  $\mathbf{z}_1 = [+1, -1, -1, +1]$ ,  $\mathbf{z}_2 = [+1, -1, +1, -1]$ ,  $\mathbf{z}_3 = [+1, +1, +1, +1]$ ,  $\mathbf{z}_4 = [-1, -1, +1, -1]$ , to perform LSH. Compute the error of the estimated cosine distance.

## 2 Programming Section [40 Points]

### 2.1 Introduction

In this coding section, we will explore a variant of approximate nearest neighbors (ANN) search for images based on locality-sensitive hashing (LSH). As discussed in class, LSH consists of a hashing function which encourages collisions between nearby points in the original space. Choosing the right hashing function is somewhat of an art and depends largely on the domain. In this assignment, we will develop the *cross-polytope hash* for image data. Roughly, the hash first projects an input sample to a lower dimensional space (by using a JL style projection) and then selects the coordinate with the highest magnitude in the resulting projected vector. The **index** of this coordinate, along with its sign, serves as a hash for the given input. One could then perform multiple such projections and stack all the individual hash values to form a **fingerprint**. The number of projections and their dimensionality is a hyperparameter that we need to tune.

## 3 Logistics and AWS

We provide the code template for this assignment is *one* Jupyter notebook. What you need to do is to follow the instructions in the notebook and implement the missing parts marked with ‘<FILL\_IN>’ or ‘# YOUR CODE HERE’.

For this homework you will be diving into using AWS’s Elastic MapReduce (EMR) systems. As such we strongly recommend that you start this homework early, as learning this is not something you will want to put off until the last moment. For setting up these services we hosted a lecture on February 15th, which guided you through the set up of everything you need to get started with AWS, we recommend going through all steps of this to familiarize yourself with the service. You can find the slides here: [AWS Guide](#)

We also plan on hosting an AWS help session recitation on Friday February 17th, which you should attend if you have any trouble setting up! If you are enrolled in the course and applied for them via the Google form, you should have \$150 AWS credits applied to your account. While we expect that all AWS based homework will cost you \$60 total this extra amount is there in case it take you more. If you use up all of your \$150 your credit card which you opened your AWS account with will be charged! Please be sure to set up usage warnings to make sure that you are not overspending. Here is a list of everything you will create which is chargeable in AWS:

1. EC2 Instances
2. EBS Volumes
3. EMR Instances
4. S3 Bucket Usage - Very small cost

When you are not working on the homework we strongly recommend you terminate/delete instances/volumes as the cost for leaving these open, even idle, can build up fast. The only thing we suggest keeping open is the S3 bucket until you are completely finished with HW4, since this costs very little. You should plan to run cheap instances while developing your code on a subset of data and only run it on something larger (and more expensive) once you have everything ready.

### 3.1 HW3 AWS Requirements and Getting lab files

For HW3 you will need to use a S3 bucket and EMR Instances. We recommend writing your code using m4.large instances using a subset of the data and using two m4.xlarge instances when you have finished your notebook and need to run on the full dataset.

You can obtain the notebooks ‘hw3.ipynb’ in the homework 3 handout .zip file. You should upload this to your Jupyterlab hosted on your AWS EMR cluster. **Be sure to download your completed notebook before you terminate your AWS cluster, otherwise you will lose all completed data.**

### 3.2 Preparing for submission

We provide several public tests via `assert` in the notebook, which will help you know if you are completing the notebook correctly. You may want to pass all those tests before submitting your homework. There is no Autograder for this homework, however in order to be given any credit for your programming section, you **MUST** upload your completed notebook to the HW3 Programming submission slot on Gradescope.

### 3.3 Deliverables

- (a) [8 points] Visualize the top 20 LSH neighbors for the test image ‘stegosaurus/image\_0043.jp’.
- (b) [8 points] Visualize the top 20 VGG neighbors for the test image ‘stegosaurus/image\_0043.jp’.
- (c) [8 points] Visualize the top 20 LSH neighbors for the test image ‘Faces/image\_0334.jpg’.
- (d) [8 points] Visualize the top 20 VGG neighbors for the test image ‘Faces/image\_0334.jpg’.
- (e) [4 points] What can you say about the neighbors returned by LSH versus those returned by VGG? Which one returns more “semantically” similar neighbors, and which one returns “pixel”-similar neighbors?
- (f) [4 points] What about the runtime of each method? Which one would be more appropriate if we had millions of images?

## 4 Collaboration Questions

1. (a) Did you receive any help whatsoever from anyone in solving this assignment?  
  
(b) If you answered ‘yes’, give full details (e.g. “Jane Doe explained to me what is asked in Question 3.4”)
  
2. (a) Did you give any help whatsoever to anyone in solving this assignment?  
  
(b) If you answered ‘yes’, give full details (e.g. “I pointed Joe Smith to section 2.3 since he didn’t know how to proceed with Question 2”)
  
3. (a) Did you find or come across code that implements any part of this assignment?  
  
(b) If you answered ‘yes’, give full details (book & page, URL & location within the page, etc.).