My Notes Important Concepts worth keeping 有函式庫可可以

```
程式碼:
```

```
To Pal (Tut str: String): booken

a Queue. create Dueue();
a stack. create Stack();

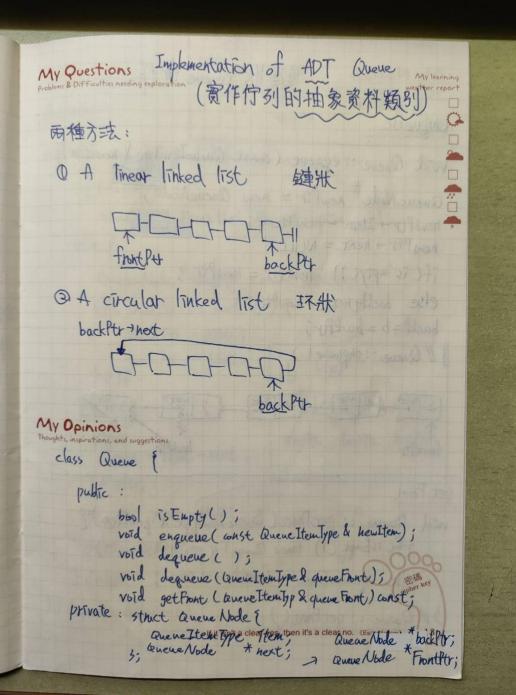
Br(the next character In str) {

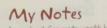
a Queue. enqueue (ch);
a stack. puch(ch);
}/for

charEqual = true; // 直展字母是方相同

while (!a Queue. is Empty() & h charEqual = = true) {

a Queue. dequeue();
a Stack. pop();
if (front! = top) charEqual = false;
}//while
```





鏈狀 Queue 實作

enqueue

void Queve :: equeve (const Queve I tem Type & new I tem) ;

Queue Node * new Ptr = new Queue Node;

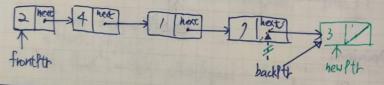
hewPtr - Item = new Item ;

hewPtr + hext = NULL;

if (is Empty()) front Ptr = new Ptr; else backlitr + next = new Ptr;

backftr = newPth;

3// Queue:: enqueue()



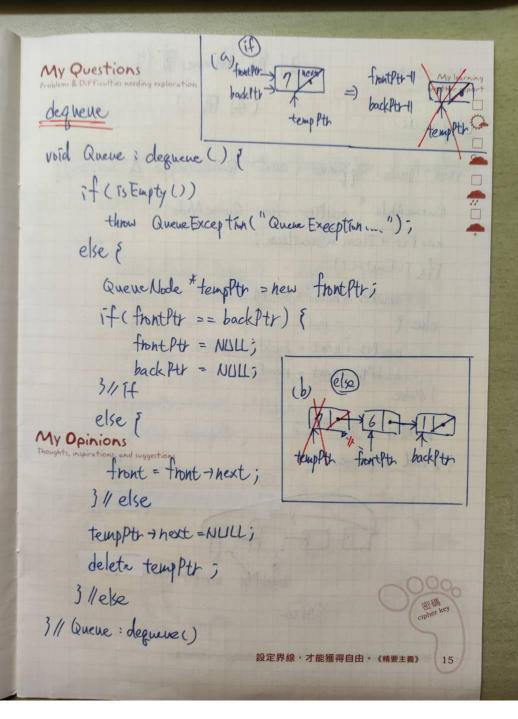
get Front

void Queue :: getfront (laveue I tem Type & queu Front) { of (is Empty()) throw Queue Exception (" "); else queue Front = front | them;

3/ Queue: getfunt()
14 If you have limits you will become limitless. (Essentialism)

VOTO

3/10



My Notes

环狀 Quele 實作 (好用!)

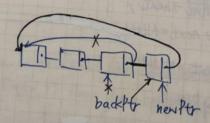
enqueve

void Queue : enqueue (anst QueueIlenType & newItem) {
QueueNode * newPter = new Queue Node;
hewPtr + Item = newItem;
If(Is Empty ())

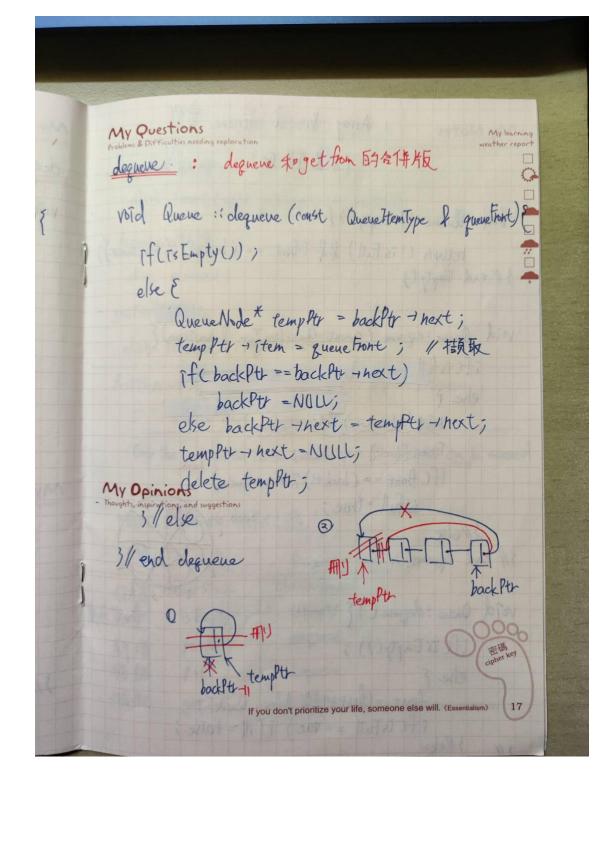
newPtr - next = new Ptr; else &

newPth - next = backPth - next; backPth + next = newPth;

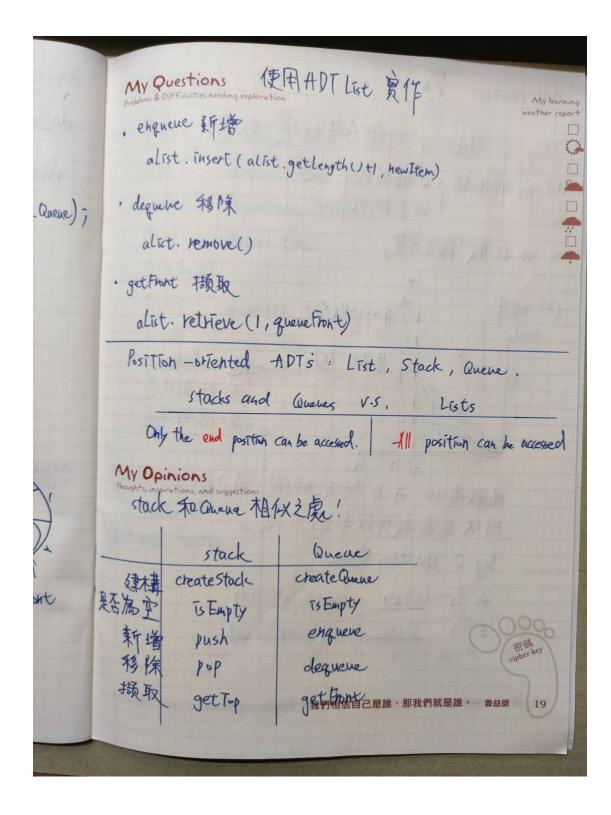
backPtr = newPtr;



16 要是你不自我管理優先緩急,自有別人為你供認



Array - based Queve \$14 My Notes (不太小董耶) al bool Quene : Is Empty()[· dequ return (! is Full) &f (front == (back+1) % Max_Quene); } / end Empty () al · getfi void ayenc: equeue (conest amenestem Type & new Item) { aLis if(isfull == true) Posit else 5 back = back+1) % Max Queue ; Items [back] = new Item; if (front == (back +1) % Max-Queue) is Full = thue; MYO back #> 3/lelse sta 3/1 end equenc() void Overe: dequeve () { 建木 是否為空 if(Is Empty()); 新! else & 18 We are what we believe we are that t) % Max - Queve ; 撷耳 if(is Full == true) is Full = false; 3/leke

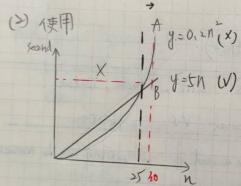


My Notes 17/2
Important Concepts worth berains

Big O Notation 判斷方法好壞(效率)

(1) 判断兩種 5 @ 時間效率

以比較「颗著差異」



隨 n 增加 A: 指數成長 B: 線性成長

比較 hin A.B 所花時間: A>B,所以 B為較有效率的方法

Brg O Notation 总注:

B is O(n) - order n (v)

20 精要主義者產出更多,其祕訣在於掃除更多,而非做更多。《精要主義》

Proble

(3

My

growth

An Feed

