

# Computer Networks - Xarxes de Computadors

#### **Outline**

- General information
- Unit 1. Introduction
- Unit 2. IP Networks
- Unit 3. LANsUnit 4. TCP
- Unit 5. Network applications

Based on: https://studies.ac.upc.edu/FIB/grau/XC/#slides



#### Lecturer

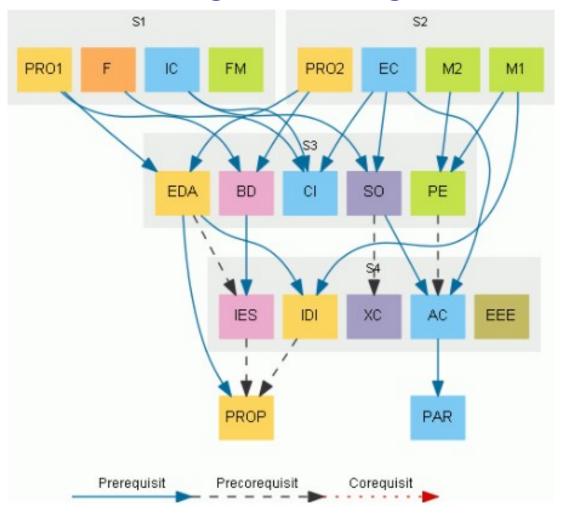
- Roger Baig Viñas roger.baig+xc@upc.edu
- Office hours (*consultes*): on demand, requests via e-mail (D6-105)

#### **Teaching resources**

- Public websites:
  - https://studies.ac.upc.edu/FIB/grau/XC/
    - Term's specific info: Read it carefully
  - https://www.fib.upc.edu/en/studies/bachelors-degrees/bachelor-degree-informatics-engineering/curriculum/syllabus/XC
    - General info
- FIB's intranet (*Racó*): https://raco.fib.upc.edu/
  - Subject: Several notifications during the semester (labs, exams, etc.)
  - Lectures: One single notification updated frequently
    - Study guidelines
    - Slides
      - » Consolidated slides will include \_final in the file name
      - » Public and possibly outdated https://people.ac.upc.edu/rbaig/XC/index.html

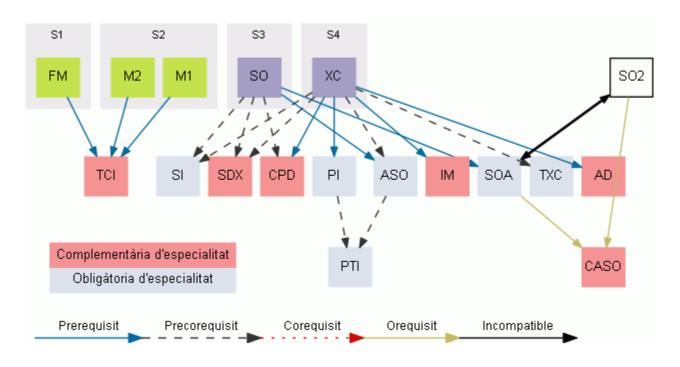


# **GEI - Assignatures obligatòries**





# GEI – Especialitat Tecnologies de la informació





#### **Course Organization**

- 2+1h lectures/week: theory + problems
  - Print the problems manual (available in the *Racó*)
  - Try to do the problems beforehand
    - Usually the problems solved will be announced in the previous lectures
  - Tracking problems (exercicis de seguiment)
    - Exercises proposed during the lectures
       One per unit, except two for Unit 2
    - Must be delivered within the next 48h
    - Can only be delivered through the Racó
      - » Where: "Pràctiques" left frame
      - » Accepted file formats: plain text (.txt), PDF, PNG, JPG
      - » Mandatory file names: Surename1\_Surename2\_Name\_seguiment{1,2,3,4,5}.ext
  - Find textbooks and related links at the subject's web page



#### **Course Organization (cont.)**

- Laboratory 7 sessions of 2h on selected weeks
  - Schedule: see subject's website (starts at 4<sup>th</sup> week)
  - Bring the manual printed in paper (@ Repography & racó)
  - Study and prepare sessions before hand
  - Submit the report at the beginning of the session
  - Minicontrol held at the end of each session (if reported submitted)



#### **Evaluation**

$$NF = 0.30 * NL + 0.70 * NT$$

- NF = Final grade
- NL = Laboratory = 0.5 \* CL + 0.5 \* EL
  - CL = Minicontrols average grade
  - EL = Laboratory final exam grade
- NT = Theory grade = 0.3 \* max(C1, EF) + 0.7 \* EF
  - C1 = Partial exam grade
  - EF = Final exam grade
    - IP's problem grade: max(C1, EF\_IP)  $\Rightarrow$  can be skipped ( $\Rightarrow$  C1 grade)

# **Bonus** (only if NF>1)

NFinc = NF + 
$$\max(0, \min(1, (NF-5)/2) * B)$$

- NFinc = Final grade incremented
- B = Tracking problems delivery rate out of 4



#### Work plan

- 6 credits ECTS (152 hours)
  - Lectures: ~39 hours
  - Laboratories: ~14 hours
  - **Self-study**: 100 hours  $\Rightarrow \sim 1.5$  of self-study hours per 1 lecture/lab hour!!
- Lectures: 13 weeks in total
  - Unit 1: Introduction 1 week
  - Unit 2: IP networks 5 weeks
  - Unit 3: LANs 2 weeks
  - Unit 4: TCP (transport) 3 weeks
  - Unit 5: Applications 2 weeks
- Exams
  - Partial (Control) 08/04/2024 10:30-12:30 Units: 1, 2
  - Final (*Examen final*) 19/06/2024 15:00-17:45 Units: all
  - Only non-programmable calculators allowed



# Work plan

- Missed with a justified cause:
  - Laboratories
    - Contact the subject's coordinator in advance (llorenc.cerda@upc.edu) to attend another group
  - Partial exam
    - The second chance exam is the final exam
  - Final exam
    - A second chance exam will be arranged



#### **Objectives**

- Identify main network functions at each level
- Identify client-server applications and associated ports
- Predict protocol operation and messages for web, e-mail, DNS apps
- Interpret documents (HTML)
- Interpret IP header fields, IP fragmentation, auxiliary protocols ARP, ICMP
- Interpret and deduce routing table content, predict RIP routing protocol behaviour and messages
- Design IP network address allocation, public and private addresses and NAT
- Design the basic configuration of a firewall (NAT, access lists and tunnels).
- Differentiate TCP and UDP and interpret header TCP segments and UDP datagrams

Source: https://www.fib.upc.edu/en/studies/bachelors-degrees/bachelor-degree-informatics-engineering/curriculum/syllabus/XC



# **Objectives (cont.)**

- Create time diagrams to model protocol behaviour in TCP
- Predict TCP flow, congestion control, window (buffers), segment transfer, congestion window, slow-start, congestion-avoidance algorithms
- Estimate the effective traffic rate for a TCP connection in different conditions (lags, link transmission speeds, segment losses, etc.)
- Represent time diagrams representing MAC protocols for the local area networks studied
- Determine the active flow control in local area network, traffic distribution in a topology (hubs, switches, routers), distinguish collision and broadcast domains, configure VLANs/ trunks and its network topology
- Identify bottlenecks in a local area network and calculate the effective flow rate for different traffic conditions.

Source: https://www.fib.upc.edu/en/studies/bachelors-degrees/bachelor-degree-informatics-engineering/curriculum/syllabus/XC



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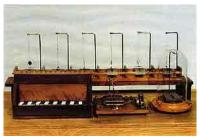
#### **Outline**

- Brief history of Computer Networks and Internet
- Introduction to the Internet
- Standardization Organizations and OSI Reference Model
- Client-Server Paradigm



# **Unit 1: Introduction Brief history of Computer Networks**

- 1830: Telegraph
- 1866: First transatlantic telegraph cable
- 1875: Alexander Graham Bell invented the telephone
- 1951: First commercial computer
- 1960: Concept of Packet-Switching.
- 1960s: ARPANET project, origins of the Internet.
- 1972: First International and commercial Packet-Switching Network, X.25.
- 1990s: The Internet is opened to the general public.



Pavel Shilling Telegraph, 1832.



Major Telegraph Lines, 1891.



UNIVAC: First commercial computer, 1951
Source: wikipedia



New York Telephone Cabling, 1888



Telephone Central Office in London, 1926

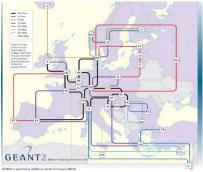


Today's Networking Equipment.



# **Brief History of the Internet**

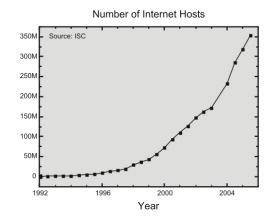
- 1966: Defense Advanced Research Projects Agency (DARPA). ARPANET project.
- ARPANET connected Universities, research labs and military centers. Military portion separated in 1983.
- 1970s: End-to-end reliability was moved to hosts, developing TCP/IP. TCP/IP was ported to UNIX Berkeley distribution, BSD.
- 1990s: The Internet is opened to commerce and the general public by the Internet Service Providers, ISP.



http://www.geant2.net



http://www.rediris.es



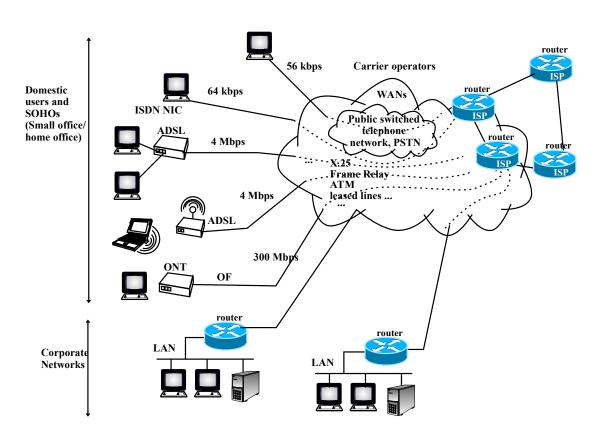


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- Host
- Access Network
- Local Area Network (LAN)
- Wide Area Network (WAN)
- Telephone company, telco, or carrier.
- Router
- Line Bitrate
- Bits per second, bps.

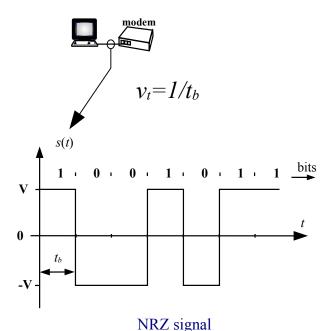




#### **Bitrate**

 $t_b$  is the transmission time of 1 bit.

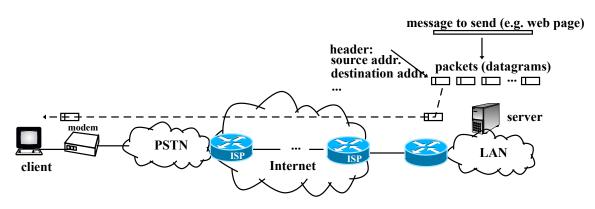
- $v_t = 1/t_b$  is the line bitrate in bits per second (bps)
- typical bitrate prefixes:
  - k, kilo: 10<sup>3</sup>
  - M, Mega: 10<sup>6</sup>
  - **G**, Giga: 10<sup>9</sup>
  - T, Tera: 10<sup>12</sup>
  - P, Peta: 10<sup>15</sup>
- Examples:
  - Public Switched Telephone Network (PSTN) modem: 56 kbps
  - ADSL: 4 Mbps
  - LAN Ethernet: 10 Mbps, 100 Mbps, 1Gbps, 10 Gbps.
  - Carrier lines E3: 34 Mbps, OC-192: 9,9 Gpbs, ...





## **Types of Switching**

- Circuit-switching, e.g. PSTN (Public Switched Telephone Network)
- Packet-switching:
  - Virtual Circuit, e.g. X.25, ATM (Asynchronous Transfer Mode).
  - Datagram: Internet.



Datagram packet- switching



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#### **Standardization Bodies**

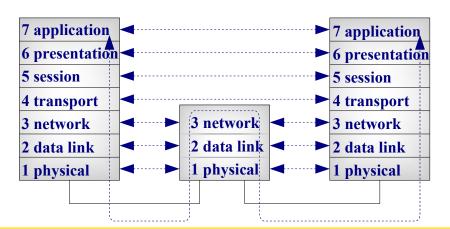
- International Telecommunication Union, ITU: WAN standards. http://www.itu.org/.
- International Organization for Standardization, ISO: Industrial standards. http://www.iso.org/.
- Institute of Electrical and Electronics Engineers, IEEE: LAN standards. http://www.ieee.org/.
- European Telecommunications Standards Institute, ETSI: Mobile phone standards (GSM). http://www.etsi.org/.
- Electronic Industries Alliance, EIA: Cabling standards. http://www.eia.org/.
- Internet Engineering Task Force, IETF: Internet standards. http://www.ietf.org. Standardization proposals are done through *Request For Comments*, RFCs. They are mirrored around the world, e.g. http://www.rfc-editor.org
- World Wide Web Consortium (W3C). http://www.w3.org



# **Unit 1: Introduction** ISO Open Systems Interconnection (OSI) Reference Model

- Layers or Levels: Physical or Layer 1 (L1), ...
- Peer layers
  - communicate using a protocol.
  - exchange *Protocol Data Unit* (PDU), which consists of a *header* and payload.
- Protocols from different layers are independent.
- Layer *i* offers services (e.g. send a datagram to a given address) to layer *i*+1: Service Access Points (SAP). Brief description of Layers:

**Terminal node ✓ → Intermediate node ✓ → Terminal node** 

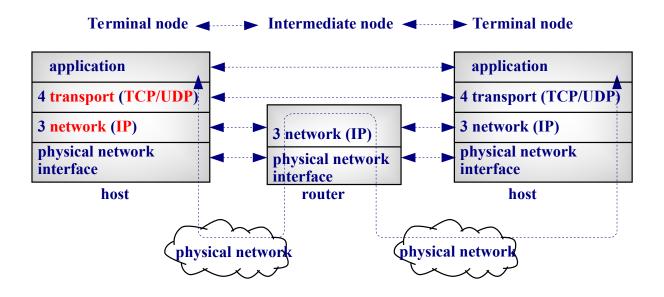


- 7. Application: Processes using network services (web, email...)
- 6. Presentation: Encoding of text, numbers...
- 5. Session: "Login" type service.
- 4. Transport: End to end data transfer.
- 3. Network: Routing.
- 2. Data link: Structured transport of bits.
- 1. Physical: Electric and mechanical.
- \*Internet jargon: Layer 8: the user.



#### **TCP/IP Architecture**

- No RFC specifies the TCP/IP model.
- Networking literature usually identifies the layer model:

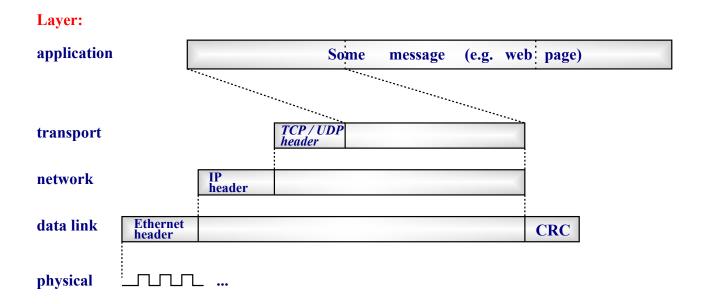


Physical network (Internet jargon): Any network that transport datagrams (not the OSI physical layer!)



# **Segmentation & Encapsulation**

Each layer adds/remove the PDU header.





#### **PDU** names

TCP/IP				
Layer	TCP	UDP		
Transport	Segment	Datagram		
Network	Packet*			
Data link	Frame			

<sup>\*</sup> Sometimes also called datagram

OSI	
Layer	
4 Transport	Segment
3 Network	Packet
2 Data link	Frame
1 Physical	Symbol (bit)



#### **Example**

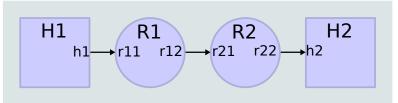
- Discussion: Just Ethernet header contents are modified over links (hops)
- Description
  - Host 1 is connected to Router 1, Router 1 to Router 2, and Router 2 to Host2
  - Host 1 sends a test packet (ICMP Echo Request) to Host 2 over the network
  - Host 2 replies with 1 packet (ICMP Echo Reply)
  - Physical network interfaces identifiers (Ethernet addresses)
    - Host 1: h1
    - Router 1: r11 and r12
    - Router 2: r21 and r22
    - Host 2: hi
  - Network identifiers (IP addresses)
    - Host 1: H1
    - Host 2: H2
- Question
  - Describe the evolution of the Ethernet headers and IP headers of both packets



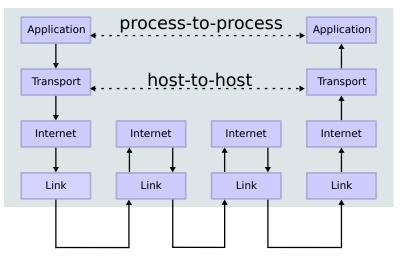
# **Example (cont.)**

Solution

**Network Topology** 



#### **Data Flow**



H1	1 ICMP ECHO
R1	1 ICMP ECHO REQUEST 6
R2	2 5/ 2 1/ICMP ECHO
H2	3 4 REPLY

		Header			
Link	Packe t	Ethernet		IP	
		src	dst	src	dst
H1-R1	1	h1	r11	H1	H2
R1-R2	2	r12	r21	H1	H2
R2-H2	3	r22	h2	H1	H2
H2-R2	4	h2	r22	H2	H1
R2-R1	5	r21	r12	H2	H1
R1-H1	6	r11	h1	H2	H1



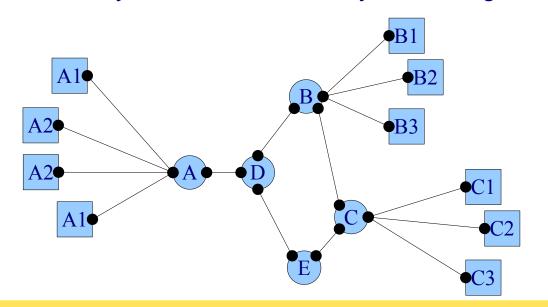
#### **Analogy: The postal service**

Network topology

Each postal office is only connected to:

- A few other postal offices
- The postal offices are connected through dedicated mailboxes (point-to-point)
- The households they serve through a single common mailbox (<u>point-to-multi-point</u>)

Any household can reach any other through at least one multi-hop path



Al Household (host)

A Postal office (router)

Mailbox (NIC)

Link



# **Analogy: The postal service (cont.)**

- Application layer We want to share a 87 pages long book with a friend
  - the book is the <u>message</u> (it is what only matters to our friend and us)
  - The only mean for sharing it is through the postal <u>service</u>
- <u>Transport layer</u> The <u>receiver</u> (our friend) limits the envelopes' capacity to 10 pages (larger envelops do not fit in her mailbox)
  - We have to split the book into several <u>segments</u>
  - The message <u>segmentation</u> and <u>reassembly</u> must be done according to the same <u>protocol</u> e.g.:
    - The first page is a special page with the number of the first page and the amount of pages of the segment. The rest of the pages are in sequential order.
  - Thus,
    - <u>Header</u>: 1 page
    - Payload: up to 9 pages
    - Segments required: 10 (9 of 9 pages each + 1 of 6 pages = 87 pages)



# **Analogy: The postal service (cont.)**

- <u>Network layer</u> Each house (<u>host</u>) has a mailbox with a <u>unique (postal/IP)</u> address
  - For each envelope:
    - We must write (won't change):
      - » Our postal address  $\rightarrow$  the source IP
      - » Our friend's postal address  $\rightarrow$  the <u>destination IP</u>
    - Put it our mailbox
  - Mailboxes (<u>routing</u>):
    - Push to the transport layer all the pages of letter meant for them (i.e. their IP as destination IP)
    - Selects the best next hop for the rest of letters and put them in the corresponding <u>queue</u> (the explanation follows in the Link layer)
  - <u>PDU</u>; : the letter as a whole
  - Header: the envelope
  - Payload: the wad of papers (i.e. up to 10 per envelope)



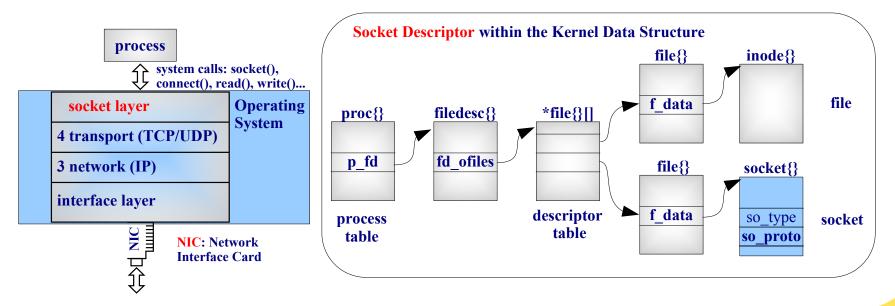
# **Analogy: The postal service (cont.)**

- Link layer
  - Mailboxes (the <u>network interface cards</u>)
    - Only one link: house closest mail offices
    - Dispatch letters encapsulating them in envelopes with the following information:
      - » Identification code of the mailbox of origin → the source MAC address
      - » Identification code of the mailbox of destination → the destination MAC address
    - Unencapsulate letters on reception (i.e the dst MAC address is theirs)
  - Mail offices (the <u>routers</u>)
    - Two links at least, one of them to a neighboring mail office
    - Routing: must decide to which of the links must enqueue each letter
    - Use the same protocol as the mail boxes
      - » Same code syntax (MAC addresses)
      - » Unencapsule all letters on reception / encapsulate all letters before sending
      - Can only send to and receive from the neighboring mail offices or their assigned houses
    - Note: here the analogy breaks a bit because in the postal service usually many envelopes are encapsulated in a single bag but in computer networks packets are encapsulated individually



# **TCP/IP Implementation**

- TCP/IP networking code is part of the Operating System kernel.
- Socket interface: Is the Unix networking interface for the processes. It was first implemented in Berkeley Software Distribution, BSD.
- The *socket* system call creates a *socket descriptor* used to store all information associated with a network connection, similarly as an inode descriptor for a file.



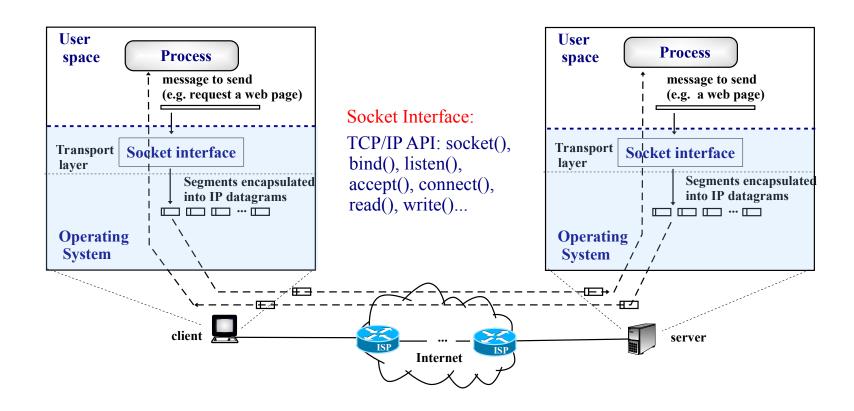


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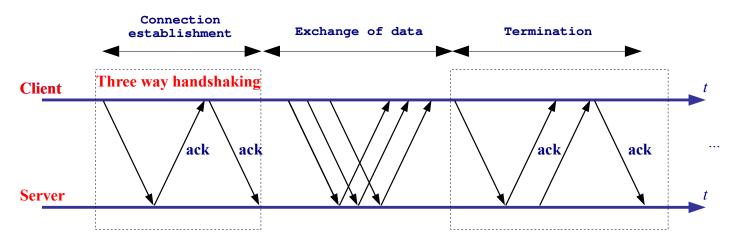
# Client Server Paradigm: Processes, messages, sockets segments and IP datagrams





# **Client Server Paradigm: The Internet Transport Layer**

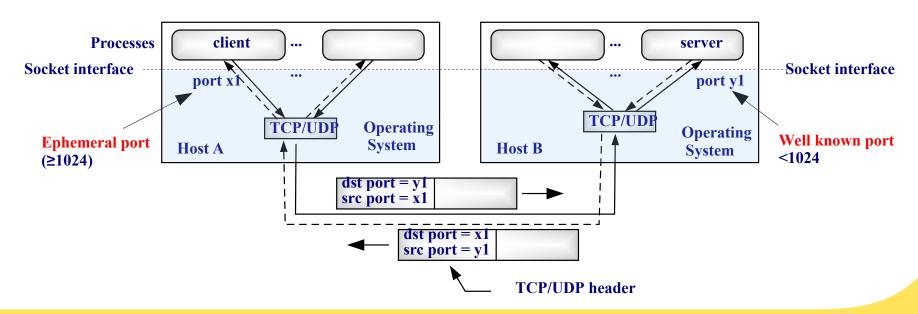
- Two protocols are used at the TCP/IP transport layer: User Datagram Protocol (UDP) and Transmission Control Protocol (TCP).
- UDP offers a *datagram service* (non reliable). It is connectionless.
- TCP offers a reliable service (correct segments are acknowledged, ack, lost segments are retransmitted). It is connection oriented (covered in detail in Unit 3).
- TCP connection:





#### **Client Server Paradigm**

- How connection is established among processes?
- The client always initiates the connection towards a known IP address, in the IP header, and a *well known* port (< 1024), in the TCP/UDP header.
- Well known ports are standardized by IANA in RFC-1700 (Assigned Numbers). In a Unix machine can be found in /etc/services.
- The server is a *daemon* waiting for client requests.





# Client Server Paradigm – UNIX /etc/services File

• Enables server and client programs to convert service names to well known ports.

```
linux> cat /etc/services
# Network services, Internet style
# Note that it is presently the policy of IANA to assign a single well-known
# port number for both TCP and UDP; hence, most entries here have two entries
# even if the protocol doesn't support UDP operations.
# This list could be found on:
             http://www.iana.org/assignments/port-numbers
# WELL KNOWN PORT NUMBERS
# The Well Known Ports are assigned by the IANA and on most systems can
# only be used by system (or root) processes or by programs executed by
# privileged users.
 Keyword Decimal Description
            7/tcp Echo
echo
            7/udp Echo
echo
          9/tcp # Discard
discard
           9/udp # Discard
discard
davtime
           13/tcp # Daytime (RFC 867)
daytime
           13/udp \# Daytime (RFC 867)
chargen
           19/tcp # Character Generator
chargen
           19/udp # Character Generator
           20/tcp # File Transfer [Default Data]
ftp-data
ftp-data
            20/udp # File Transfer [Default Data]
ftp
            21/tcp # File Transfer [Control]
            22/tcp # SSH Remote Login Protocol
ssh
            22/udp # SSH Remote Login Protocol
ssh
telnet
            23/tcp # Telnet
telnet
            23/udp # Telnet
```



# **Client Server Paradigm – Network applications**

- Remote commands
  - telnet
  - ssh
- Exchange of documents
  - ftp, sftp
  - peer-to-peer
- Web based applications
- Email
- Network management
- Real time
  - Voice over IP
  - Video streaming
- ...



# Exercici resolt: 2020t-c1-sol.pdf

Duració: 1h 30 minuts. El test es recollirà en 25 minuts.  $\longrightarrow \sim 8$  preguntes

**Test** (3,5 punts). Les preguntes valen la mitat si hi ha un error i 0 si hi ha més d'un error a la resposta.

- 1. El temps de transmissió d'un paquet de 1500 octets a 10 Mbps és 1,2 ms. En un enllaç determinat, el temps de propagació extrem a extrem entre un client i un servidor és d'1 ms. En aquest cas, el retard total extrem a extrem quan no hi ha cap node intermedi és de 2,2 ms.
- Si afegim tres routers entre el client i el servidor:
- ☐ El retard mínim extrem a extrem serà 2,2 ms.
- ☐ El retard extrem a extrem serà com a màxim 6,6 ms.
- ☐ El retard mínim extrem a extrem serà 5,8 ms.
- □ El retard mínim extrem a extrem serà 4,6 ms.

temps transmissió = 1500 octets 
$$\frac{8b}{1 \text{ octet}} \frac{1 \text{ Mb}}{1000000 \text{ b}} \frac{1}{10 \frac{\text{Mb}}{\text{s}}} = 1,2 \text{ ms}$$

retard mínim extrem a extrem = 4\*1,2 ms + 1 ms = 5,8 ms



# Exercici resolt: 2021t-c1-sol.pdf

- **2.** Dos dispositius estan connectats a través d'un router. Suposem que el temps de propagació extrem a extrem és zero, que el router no afegeix retard a les cues i que la velocitat de transmissió dels enllaços és 10 Mbps.
- ☐ Si el paquet té 1400 octets (bytes) el temps de transmissió del paquet és 0'14ms
- ☐ Si el paquet té 1400 octets (bytes) el temps de transmissió del paquet és 1'12ms.
- Si el paquet té 1400 octets (bytes) el temps total fins que ha arribat a l'altre extrem és 2'24ms.
- Si es transmeten dos paquets de 700 octets (bytes) el temps total fins que el segon paquet arriba a l'altre extrem és 1'68ms.

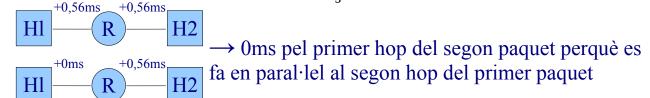
1 pk de 1400 octets

temps transmissió = 1400 octets 
$$\frac{8b}{1 \text{ octet}} \frac{1 \text{ Mb}}{1000000 b} \frac{1}{10 \frac{\text{Mb}}{\text{s}}} = 1,12 \text{ ms}$$

 $temps total \ extrem \ a \ extrem = 2 * 1,12 \ ms = 2,24 \ ms$ 

2 pk de 700 octets

temps transmissió = 700 octets 
$$\frac{8b}{1 \text{ octet}} \frac{1 \text{ Mb}}{1000000 \text{ b}} \frac{1}{10 \frac{\text{Mb}}{\text{s}}} = 0,56 \text{ ms}$$



temps total extrem a extrem = 3 \* 0.56 ms = 1.68 ms



2 El madal de vafavància ICO definais. 7 nivellas fícia, cullas de dadas, varsas transport, acceió invacantació i anlicació

# Exercicis resolts: 2021t-c1-sol.pdf

<b>3.</b> El model de referencia 150 defineix 7 nivelis. Ifsic, enliaç de dades, xarxa, transport, sessio, presentació i aplicació.				
	Tots els dispositius d'usuari i els routers de la xarxa gestionen (implementen) els 7 nivells.			
	El model de referència TCP/IP agrupa els nivells de sessió, presentació i aplicació en un únic nivell d'aplicació.			
	Tots els routers gestionen els nivells físic, enllaç de dades, xarxa i transport.			
	El nivell de transport només el gestionen els dispositius d'usuari ("hosts").			
6. Sobre el model de comunicació client-servidor.				
	Un host pot actuar a la vegada com a client i com a servidor.			
	Els paquets d'una comunicació entre processos client i servidor s'identifiquen amb les adreces IP origen i			
	destinació, els ports de client i de servidor, i el protocol.			
	Un dispositiu pot establir moltes comunicacions com a client amb el mateix servidor i protocol.			
	Un dispositiu amb una única adreça IP pot mantenir simultàniament moltes comunicacions client-servidor amb			
	molts servidors diferents.			