Design And Analysis of Algorithms - 344: Homework #1

Due on April 5, 2017 at 11:59pm

Professor Bahman Kalantari Section #1

Douglas Rudolph

Bellman Ford Algorithm

Graph $G = (V, E)$ is directed $w(u, v) \epsilon R$.

The goal is to find the shortest path from s to all other reachable verticies. In order for Bellman Ford to work, there cannot be a *Negative Cycle*. A *Negative Cycle* is a cycle where the sum of the edge weights are is less than zero.