

Fibonacci Series

0	1	2	3	4	5	6	7
0	1	1	2	3	5	8	13

Recursion

Base

case

condition

Recursive

function

call

$\downarrow T(n)$

①

fib(n) ↳

Base case
condition

if ($n = 1$) ↳

C —

return n;

↳

②

else ↳ (Recursive function

call

return fib(n-1) + fib(n-2);

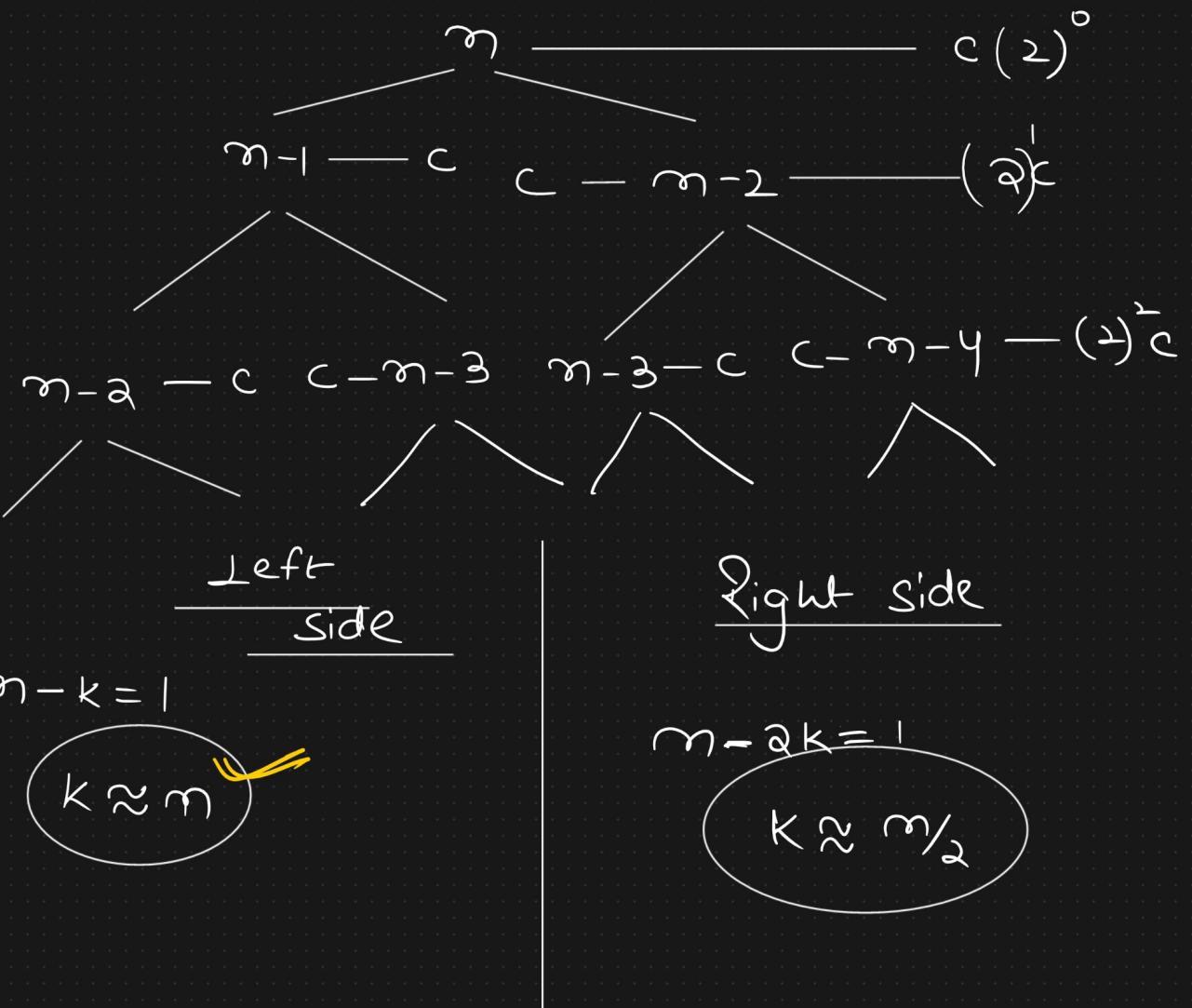
↳ $T(n-1) + T(n-2)$

↳

Recurrence Relation

\Rightarrow Recursive Tree

$$T(n) = \underbrace{T(n-1)}_{c} + \underbrace{T(n-2)}_{c} + c \quad \text{Method}$$



$$c(2^0 + 2^1 + 2^2 + \dots + 2^K) \xrightarrow[\underline{\sigma=2}]{} \text{GP series}$$

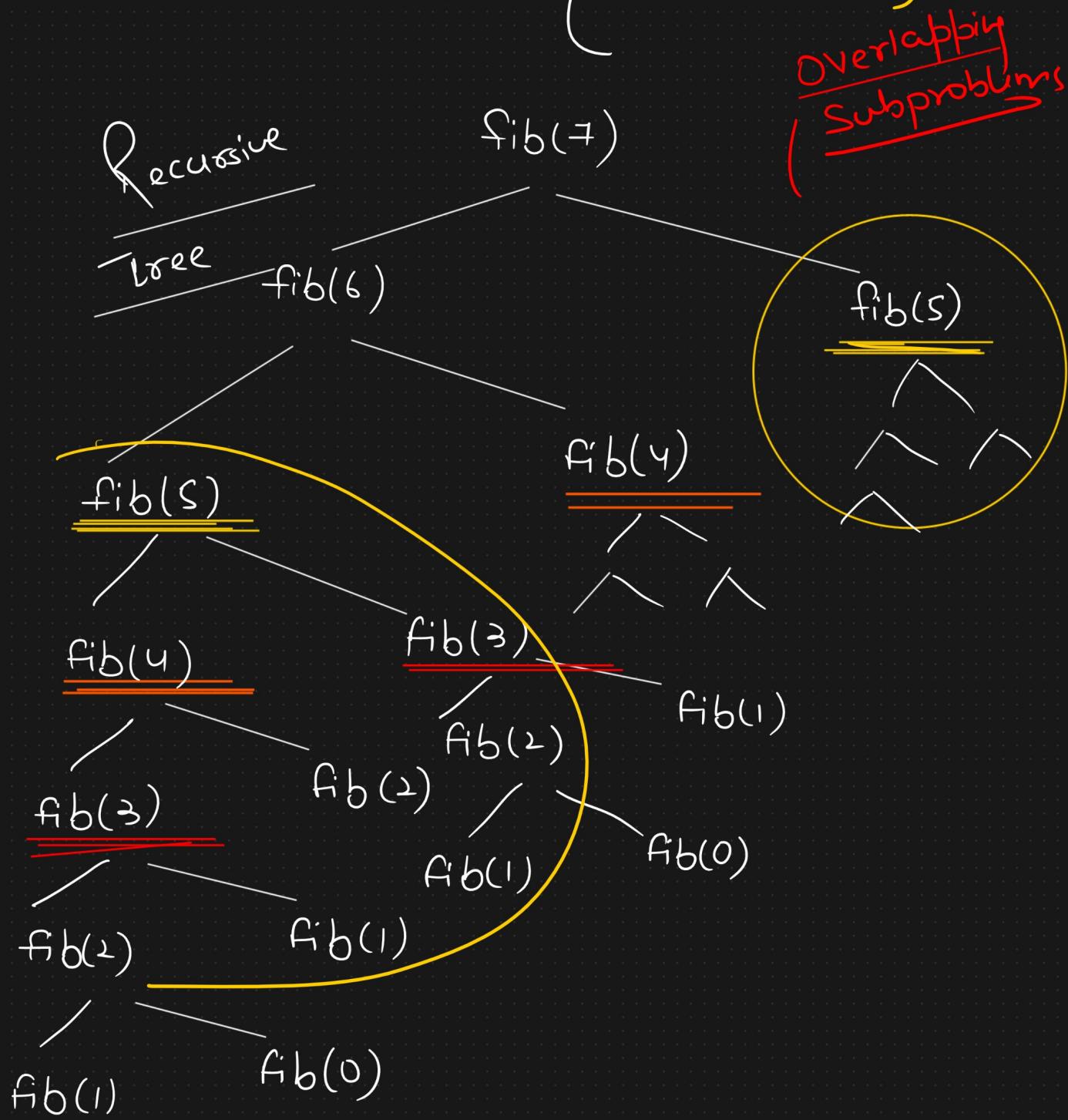
$$\frac{c(2^0 + 2^1 + 2^2 + \dots + 2^n)}{\underline{\sigma-1}} \xrightarrow{\text{Sum} = \frac{a(\sigma^n - 1)}{\sigma - 1}} \frac{c(2^n - 1)}{1} \approx O(2^n)$$

Reason ??

Exponential

time

complexity



$n \rightarrow$ very very large

$\sum 3^S$

↓

Time Limit Exceeded

① choices, Overlapping



Subproblem

↓

Re-computation of
Same function call
internally

② Optimization

(Maxima/Minima)

Dynamic Programming

① Memoization → Recursion

↓

Store the result of
every unique
function call

Exponential time
complexity

$n=7$

$\text{fib}(7)$



$\text{fib}(6)$



$\text{fib}(5)$



$\text{fib}(4)$



$\text{fib}(3)$



$\text{fib}(2)$



$\text{fib}(1)$

$\text{fib}(7)$



$\text{fib}(6)$



$\text{fib}(5)$



$\text{fib}(4)$



$\text{fib}(3)$



$\text{fib}(2)$



$\text{fib}(1)$

$\mathcal{O}(2^n)$



$\mathcal{O}(n)$



Linear

time

complexity

$\text{fib}(n) \text{ --- value}$

Hashtable

key	value
1	1
2	1
3	2
4	3
5	5
6	8

~~$SC = O(n)$~~
~~as we~~

Tabulation
No Recursion

fib	0	1	2	3	4	5	6	7
	0	1	1	2	3	5	8	13

$\text{fibonacci}(n) \propto$ time complexity
 $\text{fib}(0) = 0$
 $\text{fib}(1) = 1$

for ($i = 2$; $i \leq n$; $i++$)

\propto

$$\text{fib}(i) = \text{fib}(i-1) + \text{fib}(i-2);$$

\downarrow

return $\text{fib}(n);$

\downarrow

~~Stack Space~~

Memoization

Tabulation

→ Recursion

No Recursion

→ Slower

Faster at

Comparable to Memo.

1 Recursion \rightarrow 2 Memoization \rightarrow 3 Tabulation

4

More optimized

approach



time complexity - $O(n)$

Space complexity - $O(1)$

0	1	2	3	4	5	S
0	1	1	2	3	5	

f ~~s~~ ~~t~~ ~~f~~ ~~t~~ ~~f~~ ~~s~~ $f = 0, s = 1$

$f = s + t$

$s = f$

$t = f + s$

return s

for ($i=2$; $i \leq n$; $i++$)

$$t = f + s$$

$$f = s$$

$$s = t$$

return S;

)

{ time complexity - $O(n)$
Space complexity - $O(1)$



No use of any

external data structure
