

gcc x86 Assembly Quick Reference ("Cheat Sheet")

Instructions

Mnemonic	Purpose	Examples
mov <i>src,dest</i>	Move data between registers, load immediate data into registers, move data between registers and memory.	mov \$4,%eax # Load constant into eax mov %eax,%ebx # Copy eax into ebx mov %ebx,123 # Copy ebx to memory address 123
push <i>src</i>	Insert a value onto the stack. Useful for passing arguments, saving registers, etc.	push %ebp
pop <i>dest</i>	Remove topmost value from the stack. Equivalent to "mov (%esp), <i>dest</i> ; add \$4,%esp"	pop %ebp
call <i>func</i>	Push the address of the next instruction and start executing func.	call print_int
ret	Pop the return program counter, and jump there. Ends a subroutine.	ret
add <i>src,dest</i>	<i>dest=dest+src</i>	add %ebx,%eax # Add ebx to eax
mul <i>src</i>	Multiply eax and <i>src</i> as unsigned integers, and put the result in eax. High 32 bits of product go into eax.	mul %ebx #Multiply eax by ebx
jmp <i>label</i>	Goto the instruction <i>label</i> :. Skips anything else in the way.	jmp post_mem mov %eax,0 # Write to NULL! post_mem: # OK here...
cmp <i>a,b</i>	Compare two values. Sets flags that are used by the conditional jumps (below). WARNING: compare is relative to	cmp \$10,%eax

Stack Frame

(example without %ebp or local variables)

Contents	off esp
caller's variables	12(%esp)
Argument 2	8(%esp)
Argument 1	4(%esp)
Caller Return Address	0(%esp)

my_sub: # Returns first argument
mov 4(%esp), %eax
ret

(example when using %ebp and two local variables)

Contents	off ebp	off esp
caller's variables	16(%ebp)	24(%esp)
Argument 2	12(%ebp)	20(%esp)
Argument 1	8(%ebp)	16(%esp)
Caller Return Address	4(%ebp)	12(%esp)
Saved ebp	0(%ebp)	8(%esp)
Local variable 1	-4(%ebp)	4(%esp)
Local variable 2	-8(%ebp)	0(%esp)

my_sub2: # Returns first argument
push %ebp # Prologue
mov %esp, %ebp
mov 8(%ebp), %eax
mov %ebp, %esp # Epilogue
pop %ebp

	last argument, so "jl" jumps if $b < a$!		ret
jl <i>label</i>	Goto <i>label</i> if previous comparison came out as less-than. Other conditionals available are: jle (\leq), jeq ($==$), jge (\geq), jg ($>$), jne (\neq), and many others.	jl loop_start # Jump if $eax < 10$	
<h2>Constants, Registers, Memory</h2> <p>Constants MUST be preceeded with "\$". "\$12" means decimal 12; "\$0xF0" is hex. "\$some_function" is the address of the first instruction of the function. WARNING: a bare "12", "0xF0", or "some_function" dereferences the expression like it was a pointer!</p> <p>Registers MUST be preceeded with "%". "%eax" means register eax.</p> <p>Memory access (use register as pointer): "(%esp)". Same as C "<code>*esp</code>".</p> <p>Memory access with offset (use register + offset as pointer): "<code>4(%esp)</code>". Same as C "<code>*(esp+4)</code>".</p> <p>Memory access with scaled index (register + another register * scale): "<code>(%eax, %ebx, 4)</code>". Same as C "<code>*(eax+ebx*4)</code>".</p>			<h2>Registers</h2> <p>%esp is the stack pointer %ebp is the stack frame pointer Return value in %eax Arguments are on the stack Free for use (no save needed): %eax, %ebx, %ecx, %edx Must be saved: %esp, %ebp, %esi, %edi</p>
<h2>Common Errors</h2> <p>Segfault on innocent-looking code.</p> <ul style="list-style-type: none"> Do you need to add "\$" in front of a constant? Did you clean up the stack properly? 			

The Intel [Software Developer's Manuals](#) are incredibly long, boring, and complete--they give all the nitty-gritty details. [Volume 1](#) lists the processor registers in Section 3.4.1. [Volume 2](#) lists all the x86 instructions in Section 3.2. [Volume 3](#) gives the performance monitoring registers. For Linux, the [System V ABI](#) gives the calling convention on page 39. Also see the Intel [hall of fame](#) for historical info. [Sandpile.org](#) has a good opcode table.

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