

HIT — Cryptography — Homework 2

September 23, 2014

Problem 1. Assuming the existence of a variable output-length pseudorandom generator, present a construction of variable-length encryption scheme, and prove that your construction has indistinguishable encryptions in the presence of an eavesdropper. Hint: the construction of secure fixed-length encryption scheme also holds here.

Problem 2. Assuming the existence of a pseudorandom function, prove that there exists an encryption scheme that has indistinguishable multiple encryptions in the presence of an eavesdropper, but is not CPA-secure. Hint: You will need to use the fact that in a CPA the adversary can choose its queries to the encryption oracle adaptively (i.e., new query may be constructed from previous queries).

Problem 3. Present formulas for decryption of all the different modes of operation for encryption: ECB, CBC, OFB, CTR. For which modes can decryption be parallelized?

Problem 4. Present a construction of a variable output-length pseudorandom generator from any pseudorandom function. Prove that your construction satisfies Definition: ‘a variable output-length pseudorandom generator’.

Problem 5. Show that the CBC mode do not yield CPA-secure encryption in the case that the IV is predictable. Hint: The messages presented by the adversary could be constructed from the predictable IV and previous queries.

Problem 6. Show that the CBC, OFB and CTR modes do not yield CCA-secure encryption schemes (regardless of F). Hint: If one bit of Ciphertext is flipped, so does one bit of Plaintext.