

第五章

5.15

```
1. void inner4(ver_ptr u, vec_ptr v, data_t *dest)
2. {
3.     long i;
4.     long length = vec_length(u);
5.     data_t *udata = get_vec_start(u);
6.     data_t *vdata = get_vec_start(v);
7.     data_t s1 = (data_t)0;
8.     data_t s2 = (data_t)0;
9.     data_t s3 = (data_t)0;
10.    data_t s4 = (data_t)0;
11.    data_t s5 = (data_t)0;
12.
13.    for (i = 0; i < length; i += 6)
14.    {
15.        s1 = s1 + udata[i] * vdata[i];
16.        s2 = s2 + udata[i + 1] * vdata[i + 1];
17.        s3 = s3 + udata[i + 2] * vdata[i + 2];
18.        s4 = s4 + udata[i + 3] * vdata[i + 3];
19.        s5 = s5 + udata[i + 4] * vdata[i + 4];
20.        s6 = s6 + udata[i + 5] * vdata[i + 5];
21.    }
22.    for (; i < length; i++)
23.    {
24.        s1 = s1 + udata[i] * vdata[i];
25.    }
26.
27.    *dest = s1 + s2 + s3 + s4 + s5 + s6;
28. }
```

可能原因：

浮点数加法的单元的 issue time 为一个周期，而 Capacity 也为 1，所以每个时钟周期最多能完成 1 个假发操作，此时 CPE 下界为 1.

5.19

瓶颈在于 $val = val + a[i]$ ，这句中存在加法数据依赖，又有 $K \geq L * C$ ，浮点数的加法中 $L = 3$ ， $C = 1$ ，故选择 $3 * 1a$

```
1. void faster_psum1a(float a[], float p[], long n)
2. {
3.     long i;
4.     float val = 0;
```

```
5.     for (i = 0; (i+2) < n; i += 3)
6.     {
7.         float tmp1 = a[i];
8.         float tmp2 = tmp1 + a[i+1];
9.         float tmp3 = tmp2 + a[i+2];
10.
11.         p[i] = var + tmp1;
12.         p[i+1] = var + tmp2;
13.         p[i+2] = var = var + tmp3;
14.     }
15.     for (; i < n; ++i)
16.     {
17.         var += a[i];
18.         p[i] = var;
19.     }
20. }
```