

CCA-Secure and Authentication Encryption

Yu Zhang

Harbin Institute of Technology

Cryptography, Autumn, 2020

1 Authenticated Encryptions

2 Deterministic Encryptions

3 Key Derivation Function

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Recall Security Against CCA

The CCA indistinguishability experiment $\text{PrivK}_{\mathcal{A}, \Pi}^{\text{cca}}(n)$:

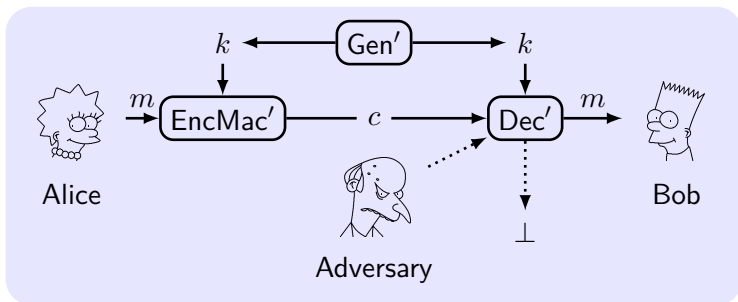
- 1 $k \leftarrow \text{Gen}(1^n)$.
- 2 \mathcal{A} is given input 1^n and oracle access $\mathcal{A}^{\text{Enc}_k(\cdot)}$ and $\mathcal{A}^{\text{Dec}_k(\cdot)}$, outputs m_0, m_1 of the same length.
- 3 a random bit $b \leftarrow \{0, 1\}$ is chosen. Then $c \leftarrow \text{Enc}_k(m_b)$ is given to \mathcal{A} .
- 4 \mathcal{A} continues to have oracle access **except for** c , outputs b' .
- 5 If $b' = b$, \mathcal{A} succeeded $\text{PrivK}_{\mathcal{A}, \Pi}^{\text{cca}} = 1$, otherwise 0.

Definition 1

Π has **indistinguishable encryptions under a CCA (CCA-secure)** if \forall PPT \mathcal{A} , \exists negl such that

$$\Pr \left[\text{PrivK}_{\mathcal{A}, \Pi}^{\text{cca}}(n) = 1 \right] \leq \frac{1}{2} + \text{negl}(n).$$

Message Transmission Scheme



- **Key-generation** algorithm outputs $k \leftarrow \text{Gen}'(1^n)$.
 $k = (k_1, k_2)$. $k_1 \leftarrow \text{Gen}_E(1^n)$, $k_2 \leftarrow \text{Gen}_M(1^n)$.
- **Message transmission** algorithm is derived from $\text{Enc}_{k_1}(\cdot)$ and $\text{Mac}_{k_2}(\cdot)$, outputs $c \leftarrow \text{EncMac}'_{k_1, k_2}(m)$.
- **Decryption** algorithm is derived from $\text{Dec}_{k_1}(\cdot)$ and $\text{Vrfy}_{k_2}(\cdot)$, outputs $m \leftarrow \text{Dec}'_{k_1, k_2}(c)$ or \perp .
- **Correctness requirement:** $\text{Dec}'_{k_1, k_2}(\text{EncMac}'_{k_1, k_2}(m)) = m$.

Defining Secure Message Transmission

The secure message transmission experiment $\text{Auth}_{\mathcal{A}, \Pi'}(n)$:

- 1 $k = (k_1, k_2) \leftarrow \text{Gen}'(1^n)$.
- 2 \mathcal{A} is given input 1^n and oracle access to EncMac'_k , and outputs $c \leftarrow \text{EncMac}'_k(m)$.
- 3 $m := \text{Dec}'_k(c)$. $\text{Auth}_{\mathcal{A}, \Pi'}(n) = 1 \iff m \neq \perp \wedge m \notin \mathcal{Q}$.

Definition 2

Π' achieves **authenticated communication** if \forall PPT \mathcal{A} , $\exists \text{negl}$ such that

$$\Pr[\text{Auth}_{\mathcal{A}, \Pi'}(n) = 1] \leq \text{negl}(n).$$

Definition 3

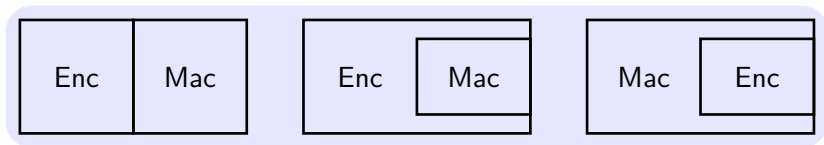
Π' is **secure Authenticated Encryption (A.E.)** if it is both CCA-secure and also achieves authenticated communication.

Does CCA-security imply A.E.? (homework)

Suppose (E, D) provides A.E. Which of the following systems provide A.E.?

- $E'_k(m) = (E_k(m), E_k(m))$ and $D'_k(c_1, c_2) = D_k(c_1)$
- $E'_k(m) = (E_k(m), 0)$ and $D'_k(c, b) = \begin{cases} D_k(c) & \text{if } b = 0 \\ \perp & \text{otherwise} \end{cases}$
- $E'_k(m) = (E_k(m), E_k(m))$ and $D'_k(c_1, c_2) = \begin{cases} D_k(c_1) & \text{if } D_k(c_1) = D_k(c_2) \\ \perp & \text{otherwise} \end{cases}$
- $E'_k(m) = (E_k(m), H(m))$ (H is a CRHF) and $D'_k(c, h) = \begin{cases} D_k(c) & \text{if } H(D_k(c)) = h \\ \perp & \text{otherwise} \end{cases}$

Combining Encryption and Authentication



- **Encrypt-and-authenticate** (e.g., SSH):

$$c \leftarrow \text{Enc}_{k_1}(m), t \leftarrow \text{Mac}_{k_2}(m).$$

- **Authenticate-then-encrypt** (e.g., SSL):

$$t \leftarrow \text{Mac}_{k_2}(m), c \leftarrow \text{Enc}_{k_1}(m \| t).$$

- **Encrypt-then-authenticate** (e.g., IPsec):

$$c \leftarrow \text{Enc}_{k_1}(m), t \leftarrow \text{Mac}_{k_2}(c).$$

Analyzing Security of Combinations

All-or-nothing: Reject any combination for which there exists even a single counterexample is insecure.

- **Encrypt-and-authenticate:** $\text{Mac}'_k(m) = (m, \text{Mac}_k(m))$.

- **Authenticate-then-encrypt:**

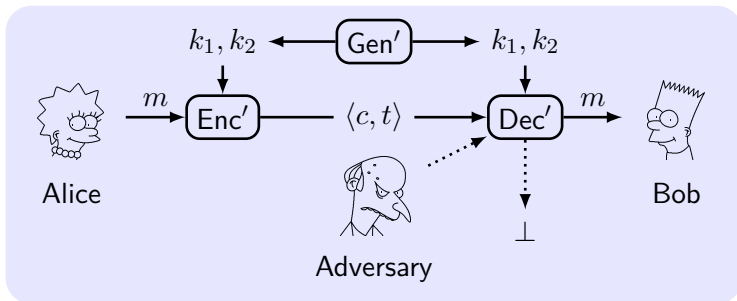
- $\text{Trans} : 0 \rightarrow 00; 1 \rightarrow 10/01$; Enc' uses CTR mode;
 $c = \text{Enc}'(\text{Trans}(m \parallel \text{Mac}(m)))$.
- Flip the first two bits of c and verify whether the ciphertext is valid. $10/01 \rightarrow 01/10 \rightarrow 1$, $00 \rightarrow 11 \rightarrow \perp$.
- If valid, the first bit of message is 1; otherwise 0.
- For any MAC, this is not CCA-secure.

- **Encrypt-then-authenticate:**

Decryption: If $\text{Vrfy}(\cdot) = 1$, then $\text{Dec}(\cdot)$; otherwise output \perp .

Constructing Authenticated Encryption Schemes

Idea: Make decryption oracle useless. AE(/CCA-secure) = CPA-then-MAC.



Construction 4

$\Pi_E = (\text{Gen}_E, \text{Enc}, \text{Dec})$, $\Pi_M = (\text{Gen}_M, \text{Mac}, \text{Vrfy})$. Π' :

- $\text{Gen}'(1^n)$: $k_1 \leftarrow \text{Gen}_E(1^n)$ and $k_2 \leftarrow \text{Gen}_M(1^n)$
- $\text{Enc}'_{k_1, k_2}(m)$: $c \leftarrow \text{Enc}_{k_1}(m)$, $t \leftarrow \text{Mac}_{k_2}(c)$ and output $\langle c, t \rangle$
- $\text{Dec}'_{k_1, k_2}(\langle c, t \rangle) = \text{Dec}_{k_1}(c)$ if $\text{Vrfy}_{k_2}(c, t) \stackrel{?}{=} 1$; otherwise \perp

Proof of AE(/CCA-Secure) Encryption Schemes

Theorem 5

If Π_E is a CPA-secure private-key encryption scheme and Π_M is a secure MAC with unique tags, then Construction Π' is AE(/CCA-secure).

Proof.

VQ: \mathcal{A} submits a “new” query to oracle Dec' and $\text{Vrfy} = 1$.

$$\Pr[\text{PrivK}_{\mathcal{A}, \Pi'}^{\text{cca}}(n) = 1] \leq \Pr[\text{VQ}] + \Pr[\text{PrivK}_{\mathcal{A}, \Pi'}^{\text{cca}}(n) = 1 \wedge \overline{\text{VQ}}]$$

We need to prove the following claims.

- 1 $\Pr[\text{VQ}]$ is negligible.
- 2 $\Pr[\text{PrivK}_{\mathcal{A}, \Pi'}^{\text{cca}}(n) = 1 \wedge \overline{\text{VQ}}] \leq \frac{1}{2} + \text{negl}(n)$.



Proof of “ $\Pr[\text{VQ}]$ is negligible”

Idea: Reduce \mathcal{A}_M (attacking Π_M with an oracle $\text{Mac}_{k_2}(\cdot)$) to \mathcal{A} .

Proof.

- \mathcal{A}_M chooses $i \leftarrow \{1, \dots, q(n)\}$ *u.a.r.*
- Run \mathcal{A} with the encryption/decryption oracles.
- If the i th decryption oracle query from \mathcal{A} uses a “new” c , output (c, t) and stop.
- $\text{Macforge}_{\mathcal{A}_M, \Pi_M}(n) = 1$ only if VQ occurs.
- \mathcal{A}_M correctly guesses i with probability $1/q(n)$.

$$\Pr[\text{Macforge}_{\mathcal{A}_M, \Pi_M}(n) = 1] \geq \Pr[\text{VQ}]/q(n).$$



Proof of “ $\Pr[\text{PrivK}_{\mathcal{A}, \Pi'}^{\text{cca}}(n) = 1 \wedge \overline{\text{VQ}}] \leq \frac{1}{2} + \text{negl}(n)$ ”

Idea: Reduce \mathcal{A}_E (attacking Π_E with an oracle $\text{Enc}_{k_1}(\cdot)$) to \mathcal{A} .

Proof.

- Run \mathcal{A} with the encryption/decryption oracles.
- Run $\text{PrivK}_{\mathcal{A}_E, \Pi_E}^{\text{cpa}}$ as $\text{PrivK}_{\mathcal{A}, \Pi'}^{\text{cca}}$.
- \mathcal{A}_E outputs the same b' that is output by \mathcal{A} .
- $\Pr[\text{PrivK}_{\mathcal{A}_E, \Pi_E}^{\text{cpa}}(n) = 1 \wedge \overline{\text{VQ}}] = \Pr[\text{PrivK}_{\mathcal{A}, \Pi'}^{\text{cca}}(n) = 1 \wedge \overline{\text{VQ}}]$ unless VQ occurs.

$$\Pr[\text{PrivK}_{\mathcal{A}_E, \Pi_E}^{\text{cpa}}(n) = 1] \geq \Pr[\text{PrivK}_{\mathcal{A}, \Pi'}^{\text{cca}}(n) = 1 \wedge \overline{\text{VQ}}].$$



Theorem 6

Π_E is CPA-secure and Π_M is a secure MAC with unique tags, Π' deriving from encrypt-then-authenticate approach is secure.

GCM(Galois/Counter Mode): CTR encryption then Galois MAC. (RFC4106/4543/5647/5288 on IPsec/SSH/TLS)

EAX: CTR encryption then CMAC (Cipher-based MAC).

Proposition 7

Authenticate-then-encrypt approach is secure if Π_E is rand-CTR mode or rand-CBC mode.

CCM (Counter with CBC-MAC): CBC-MAC then CTR encryption. (802.11i, RFC3610)

OCB (Offset Codebook Mode): integrating MAC into encryption. (two times fast as CCM, EAX)

All support AEAD (A.E. with associated data): part of message is in clear, and all is authenticated

Remarks on Secure Message Transmission

- Authentication may leak the message.
- Secure message transmission implies CCA-security. The opposite direction is not necessarily true.
- Different security goals should always use different keys.
 - otherwise, the message may be leaked if $\text{Mac}_k(c) = \text{Dec}_k(c)$.
- Implementation may destroy the security proved by theory.
 - **Attack with padding oracle** (in TLS 1.0):
Dec return two types of error: padding error, MAC error.
Adv. learns last bytes if no padding error with guessed bytes.
 - **Attack non-atomic dec.** (in SSH Binary Packet Protocol):
Dec (1)decrypt length field; (2)read packets as specified by the length; (3)check MAC.
Adv. (1)send c ; (2)send l packets until “MAC error” occurs; (3)learn $l = \text{Dec}(c)$.

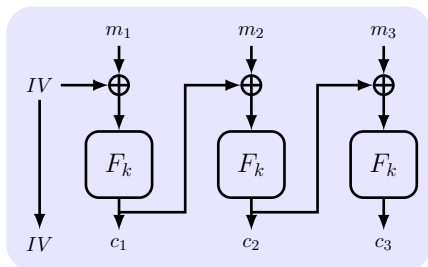
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Deterministic CPA Security

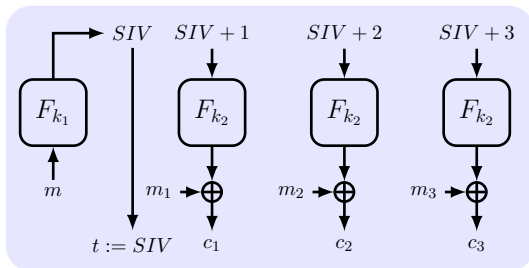
- **Application:** encrypted database index, disk encryption
- **Deterministic encryption:** the same message is encrypted to the same ciphertext under the same key.
- **Q: would it be CPA-secure?**
- **Deterministic CPA Security:** CPA-secure if *never encrypt same message twice* using same key. The pair $\langle k, m \rangle$ is unique.
- **Deterministic Authenticated Encryption (DAE)**
- **Common Mistake:** CBC/CTR with **fixed IV**.



Adversary can query (m_{q1}, m_{q2}) and get (c_{q1}, c_{q2}) ; then output PT: $IV \oplus c_{q1} \oplus m_{q2}$ and expect CT: c_{q2} .

Synthetic IV (SIV) for Det. Encryption

- **SIV** (fixed IV for the same $\langle k, m \rangle$):
PRF F , CPA-secure $\Pi : (\text{Enc}_k(r, m), \text{Dec}_k(r, s))$
 $(k_1, k_2) \leftarrow \text{Gen}$; $SIV \leftarrow F_{k_1}(m)$
 $c = \langle SIV, \text{Enc}_{k_2}(SIV, m) \rangle$.
- **DAE for free with SIV-CTR**:

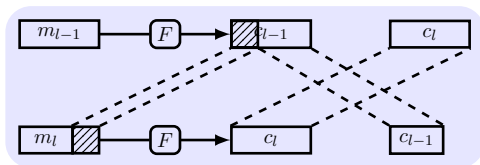


Wide Block PRP for Det. Encryption

- **Wide block PRP:** PRP with longer block length (e.g. a sector on disk) from PRP with short block length (e.g. AES).
- **PRP-based DAE:** If F is PRP, then F is also det. CPA-secure. $\text{Enc}_k(m\|0^\ell)$. In Dec, if $\neq 0^\ell$, output \perp .
- **Narrow block** may leak info. as some blocks are the same.
- **Standards:** CBC-mask-CBC (CMC) and ECB-mask-ECB (EME) in IEEE P1619.2.
- **Cost:** 2x slower than SIV due to two-pass encryption.

Tweakable Encryption

- **Encryption without expansion:** $\mathcal{M} = \mathcal{C}$ implies det. encryption without integrity (e.g., disk encryption).
- **Tweak:** like IV , different tweak for different block.
- **Trivial solution:** $k_t = F_k(t), t = 1, \dots, \ell$.
- **Tweakable block ciphers:** many PRPs from one key $\mathcal{K} \times \mathcal{T} \times \mathcal{X} \rightarrow \mathcal{X}$, \mathcal{T} is the set of tweaks.
- **XTS:** XEX(Xor-Encrypt-Xor)-based tweaked-codebook mode with ciphertext stealing. (XTS-AES, NIST SP 800-38E)
- **XEX:** To encrypt block j in sector I , $c = F_k(m \oplus x) \oplus x$, where $x = F_k(I) \otimes 2^j$ in Galois field, (I, j) is tweak.
- **Ciphertext stealing (CTS):** no padding, **no expansion**.



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Key Derivation Function (KDF)

Key Derivation Function (KDF) generates many keys from a secret source key sk .

For uniformly random sk : F is PRF, ctx is a unique string identifying application,

$$\text{KDF}(sk, ctx, l) = \langle F_{sk}(ctx\|0), F_{sk}(ctx\|1) \cdots, F_{sk}(ctx\|l) \rangle.$$

For not-uniform sk : extract-then-expand paradigm.

extract: HKDF $k \leftarrow \text{HMAC}(salt, sk)$. $salt$ is a random number.

expand: as the above.

Password-Based KDF (PBKDF)

Key stretching increases the time of testing key (with slow hash function).

Key strengthening increases the length/randomness of key (with salt).

PKCS#5 (PBKDF1): $H^{(c)}(pwd || salt)$, iterate hash function c times.

Attacker: either try the enhanced key (larger key space), or else try the initial key (longer time per key).

IV, Nonce, Counter and Salt

IV an input to a cryptographic primitive, providing randomness.

nonce a number used only once to sign a communication.

counter a sequence number used as nonce or IV.

tweak an input used only once for each block in a cipher.

salt consists of random bits, creating the input to a function.

- CCA-secure, Authenticated Encryption, Enc-then-Auth(/MAC).
- SIV, Wide Block Cipher, Tweakable Encryption, CTS.
- PBKDF, salt.