

# Recurrences

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        result *= n;
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  - **base case**

## Example 1: factorial()

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```
int factorial (int n)
{
    if (n == 0)                       $T(n) =$ 
        return 1;

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# Recurrence is algorithmic

A recurrence  $T(n)$  is *algorithmic* if, for every sufficiently large threshold constant  $n_0 > 0$ , the following properties hold:

- ①  $\forall n < n_0, T(n) = \Theta(1)$
- ②  $\forall n \geq n_0$ , every path of recursion terminates in a defined base case within a finite number of recursive invocations.

**Convention for recurrences:** *Whenever a recurrence is stated without an explicit base case, we assume that the recurrence is algorithmic.*

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void displayLine (int n)
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    if (n > 0) {
        print ('=');
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# Exercise 1

Write the recursive function and derive its recurrence relation.

- ① `sum (n)` : computes for  $\sum_{j=1}^n j$
- ② `power (x, y)` : computes for  $x^y$
- ③ `fibonacci (n)` : returns the  $n^{\text{th}}$  fibonacci number.

Fibonacci numbers are: 0, 1, 1, 2, 3, 5, 8, 13, 21, ...

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Determine the closed form or explicit formula of  $T(n)$ .

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$T(n) = 3n - 1$	$O(n)$	(9)
-----------------	--------	-----

## Exercise 2

Use the iteration method to determine the explicit formula and the *Big-Oh* of the following recurrences defined.

$$\textcircled{1} \quad T(n) = \begin{cases} a & n = 1 \\ T(n - 1) + b & n > 1 \end{cases}$$

$$\textcircled{2} \quad T(n) = \begin{cases} 5 & n = 0 \\ T(n - 1) + 2n - 1 & n > 0 \end{cases}$$

$$\textcircled{3} \quad T(n) = \begin{cases} 1 & n = 1 \\ T\left(\frac{n}{2}\right) + 1 & n > 1 \end{cases}$$

$$\textcircled{4} \quad T(n) = \begin{cases} 8 & n = 1 \\ 2T(n - 1) + 5 & n > 1 \end{cases}$$

$$\textcircled{5} \quad T(n) = \begin{cases} 1 & n = 1 \\ 2T\left(\frac{n}{2}\right) + 1 & n > 1 \end{cases}$$

# Formulas to recall : Summations

$$\sum_{j=a}^b 1 = b - a + 1$$

$$\sum_{j=0}^n j = \frac{n(n+1)}{2}$$

$$\sum_{j=0}^n j^2 = \frac{n(n+1)(2n+1)}{6}$$

$$\sum_{j=0}^n j^3 = \frac{n^2(n+1)^2}{4}$$

$$\sum_{j=0}^n x^j = \frac{x^{n+1} - 1}{x - 1}$$

# Logarithms

Notations:

$$\lg n = \log_2 n \quad \text{binary logarithm}$$

$$\ln n = \log_e n \quad \text{natural logarithm}$$

$$\lg^k n = (\lg n)^k \quad \text{exponentiation}$$

$$\lg \lg n = \lg(\lg n) \quad \text{composition}$$

$$\log n + 1 = (\log n) + 1 \neq \log(n + 1)$$

For any constant  $b > 1$ ,

$$\log_b n = \begin{cases} \text{undefined} & n \leq 0 \\ \text{strictly increasing} & n > 0 \\ \text{negative} & 0 < n < 1 \\ \text{positive} & n > 1 \\ 0 & n = 1 \end{cases}$$

# Formulas to remember : Logarithms

For all real  $a > 0, b > 0, c > 0, n$ , and logarithm bases  $\neq 1$

$$a = b^{\log_b a}$$

$$\log_c ab = \log_c a + \log_c b$$

$$\log_b a^n = n \log_b a$$

$$\log_b \frac{1}{a} = -\log_b a$$

$$\log_b a = \frac{\log_c a}{\log_c b}$$

$$\log_b a = \frac{1}{\log_a b}$$

$$a^{\log_b c} = c^{\log_b a}$$

# Reference

Antioquia, A. (2020). Recurrences. CCDSALG Lecture.

Cormen, T., Leiserson, C., Rivest, R., and Stein, C. (2022). Introduction to Algorithms, fourth edition. MIT Press.