

CSARCH Lecture Series: Floating-Point Operation and Guard Bits

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Overview

Reflect on the following question:

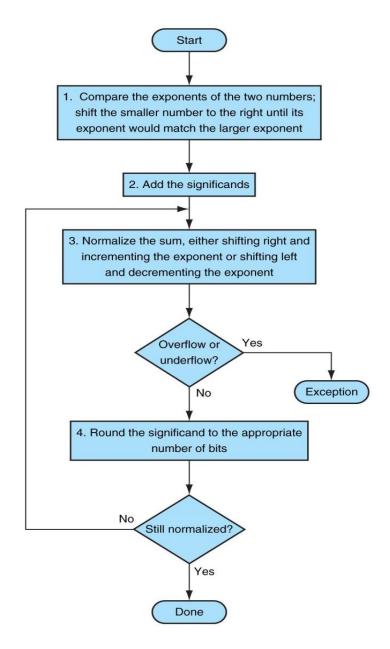
 Consider a system that can support 7 bits, perform FP addition using guard, round and sticky bits

 $> 1.0111110010_2 \times 2^5 + 1.00111111110_2 \times 2^3$

Overview

- This sub-module introduces floating-point arithmetic operations and the concepts of guard, round and sticky bits
- The objectives are as follows:
 - ✓ Describe the process of performing floating-point arithmetic operations
 - ✓ Describe the process of using guard, round and sticky bits

FP Addition Algorithm



Floating-point addition. The normal path is to execute steps 3 and 4 once, but if rounding causes the sum to be unnormalized, we must repeat step 3.

Floating-Point Addition

- Consider a system that can support 4 decimal digits
 - $> 9.999 \times 10^1 + 1.610 \times 10^{-1}$
- 1. Align decimal points

$$\triangleright$$
 9.999 × 10¹ + 0.016 × 10¹

2. Add the significands/significand

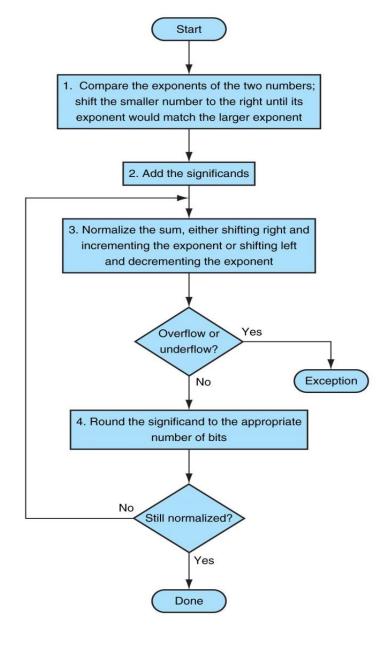
$$> 9.999 \times 10^1 + 0.016 \times 10^1 = 10.015 \times 10^1$$

3. Normalize result & check for over/underflow

$$> 1.0015 \times 10^2$$

4. Round and renormalize if necessary

$$> 1.002 \times 10^2$$



Floating-Point Addition

Consider a system that can support 4 bits:

$$> 1.000_2 \times 2^{-1} + (-1.110_2) \times 2^{-2}$$

1. Align binary points

$$> 1.000_2 \times 2^{-1} + (-0.111_2) \times 2^{-1}$$

2. Add significands

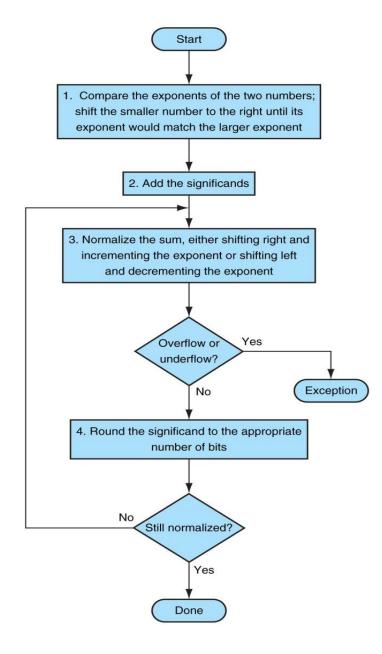
$$> 1.000_2 \times 2^{-1} + (-0.111_2) \times 2^{-1} = 0.001_2 \times 2^{-1}$$

3. Normalize result & check for over/underflow

$$\geq$$
 1.000₂ × 2⁻⁴

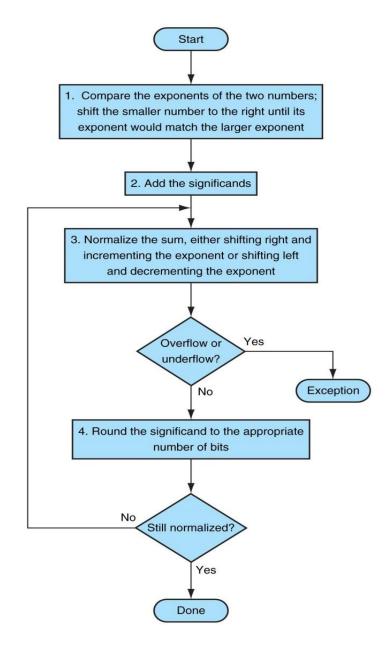
4. Round and renormalize if necessary

$$> 1.000_2 \times 2^{-4}$$
 (no change)



Floating-Point Subtraction

- Consider a system that can support 7 decimal digits
 - **▶** 123457.1467 123456.659
- 1. Align decimal points
 - \triangleright 1.234571467 x 10⁵ 1.23456659 x 10⁵
- 2. Subtract significands
 - \triangleright 1.234571 x 10⁵ 1.234567 x 10⁵ = 0.000004 × 10⁵
- 3. Normalize result & check for over/underflow
 - $> 4.0000000 \times 10^{-1}$
- 4. Round and renormalize if necessary
 - > 4.000000 \times 10⁻¹ (same)



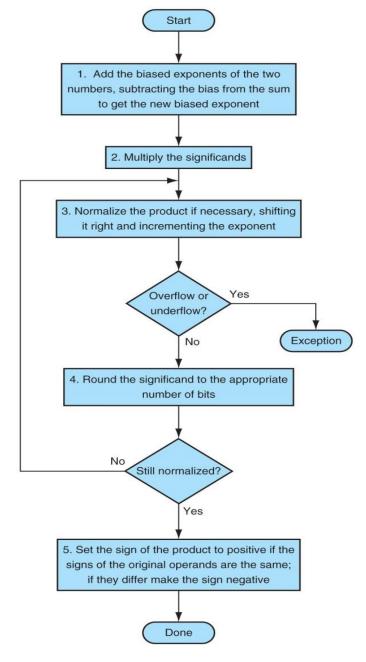
Floating-Point Addition

Reflection: Floating-point is not associative

A= 1234.567, b=45.67834, c=0.0004

Is
$$(A+B)+C = A+(B+C)$$
?

FP Multiplication Algorithm



Floating-Point Multiplication

- Consider a system that can support 4 decimal digits
 - $> 1.110 \times 10^{10} \times 9.200 \times 10^{-5}$
- 1. Add exponents
 - \triangleright New exponent = 10 + -5 = 5
- 2. Multiply significands

$$\triangleright$$
 1.110 × 9.200 = 10.212 \Rightarrow 10.212 × 10⁵

3. Normalize result & check for over/underflow

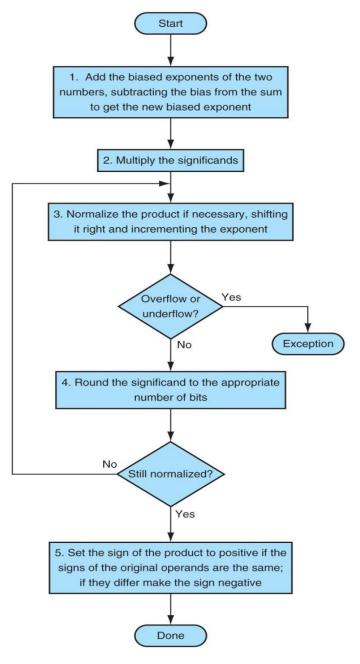
$$> 1.0212 \times 10^6$$

4. Round and renormalize if necessary

$$\geq$$
 1.021 × 10⁶

5. Determine sign of result from signs of operands

$$> +1.021 \times 10^6$$



Floating-Point Multiplication

Consider a system that can support 4 bits

$$\triangleright 1.000_2 \times 2^{-1} \times -1.110_2 \times 2^{-2}$$

1. Add exponents

$$\triangleright$$
 Unbiased: $-1 + -2 = -3$

2. Multiply significands

$$\triangleright 1.000_2 \times 1.110_2 = 1.110_2 \implies 1.110_2 \times 2^{-3}$$

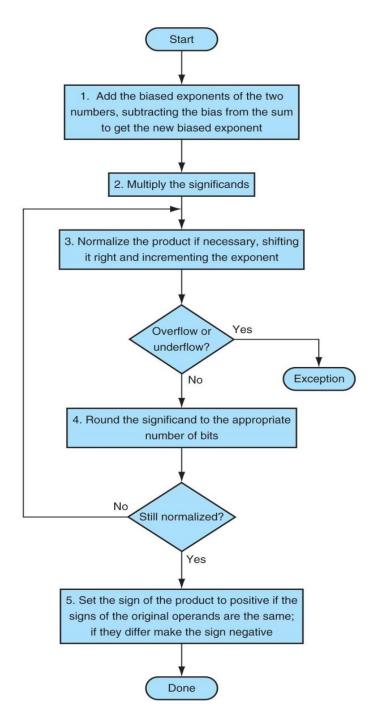
3. Normalize result & check for over/underflow

$$> 1.110_2 \times 2^{-3}$$
 (no change)

4. Round and renormalize if necessary

$$> 1.110_2 \times 2^{-3}$$
 (no change)

5. Determine sign: +operand \times -operand \Rightarrow -operand \Rightarrow -1.110₂ \times 2⁻³



Guard, Round & Sticky Bits

- Guard bit the first of the two extra bits kept on the right during intermediate calculations of floating point numbers; used to improve rounding accuracy.
- Round bit the second of the two extra bits kept on the right during intermediate calculations of floating point numbers; used to improve rounding accuracy.
- Sticky bit bit used in rounding in addition to guard and round bit that is set whenever there are nonzero bits to the right of the round bit.

Guard, Round & Sticky Digits

• For decimal, the sticky digit is always 1 if there is a non-zero in the least significant digit

Rounding with guard, round & sticky bits

- Consider a binary that can support 5 bits
 - \geq 1.00111101 x 2⁵ + 1.00111101 x 2³
- 1. Align binary points
 - ➤ Shift number with smaller exponent
 - \triangleright 1.00111101 x 2⁵ + 0.0100111101 x 2⁵
- 2. Add significands (with guard, round and sticky bits)
 - \geq 1.0011111 x 2⁵ + 0.0100111 x 2⁵ = 1.1000110 × 2⁵
- 3. Normalize result & check for over/underflow
 - $> 1.1001 \times 2^5$
- 4. Round and renormalize if necessary
 - $> 1.1001 \times 2^5$

Rounding without guard, round & sticky bits

- Consider a binary that can support 5 bits
 - \triangleright 1.00111101 x 2⁵ + 1.00111101 x 2³
- 1. Align binary points
 - ➤ Shift number with smaller exponent
 - \triangleright 1.00111101 x 2⁵ + 0.0100111101 x 2⁵
- 2. Add significands
 - \geq 1.0100 x 2⁵ + 0.0101 x 2⁵ = 1.1001 × 2⁵
- 3. Normalize result & check for over/underflow
 - $> 1.1001 \times 2^5$
- 4. Round and renormalize if necessary
 - $> 1.1001 \times 2^5$

Rounding with guard, round & sticky digits

- Consider a system that can support 7 decimal digits
 - **▶** 123457.1467 123456.659
- 1. Align decimal points
 - ➤ Shift number with smaller exponent
 - \triangleright 1.234571467 x 10⁵ 1.23456659 x 10⁵
- 2. Add significands (with guard, round and sticky digits)
 - \geq 1.234571461 x 10⁵ 1.234566590 x 10⁵ = 0.000004871 × 10⁵
- 3. Normalize result & check for over/underflow
 - $>4.871000 \times 10^{-1}$
- 4. Round and renormalize if necessary
 - \geq 4.871000 \times 10⁻¹ vs. [compare to 4.000000 \times 10⁻¹ without guard & round digits]



 $> 1.0111110010_2 \times 2^5 + 1.00111111110_2 \times 2^3$

	Base ^{Exp}
Operand 1	
Operand 2	
Sum	
Sum (normalize)	

Perform without guard, round and sticky bits



 $> 1.0111110010_2 \times 2^5 + 1.00111111110_2 \times 2^3$

		Base ^{Exp}
Operand 1	1.011111	2 ⁵
Operand 2	0.010100	2 ⁵
Sum	1.110011	2 ⁵
Sum (normalize)	1.110011	2 ⁵

Perform without guard, round and sticky bits



 $> 1.0111110010_2 \times 2^5 + 1.00111111110_2 \times 2^3$

	G	R	S	Base ^{Exp}
Operand 1				
Operand 2				
Sum				
Sum (normalize)				

Perform with guard(G), round(R) and sticky(S) bits



 $> 1.0111110010_2 \times 2^5 + 1.00111111110_2 \times 2^3$

		G	R	S	Base ^{Exp}
Operand 1	1.011111	0	0	1	2 ⁵
Operand 2	0.010011	1	1	1	2 ⁵
Sum	1.110011	0	0	0	2 ⁵
Sum (normalize)	1.110011	-	-	-	2 ⁵

Perform with guard(G), round(R) and sticky(S) bits

To recall ...

- What have we learned:
- The objective are as follows:
 - ✓ Describe the process of performing FP arithmetic operations
 - ✓ Describe the process of using guard, round and sticky bits