

CSARCH Lecture Series: Non-Restoring Division

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Overview

Reflect on the following question:

- How does Arithmetic and Logic Unit (ALU) perform division?

```
int main()
{
    int var, var1, var2;
    var = 5;
    var1 = 2;
    var2 = var / var1;
}
```



Overview

- This sub-module describes how Arithmetic Logic Unit (ALU) performs division using non-restoring division
- The objective is as follows:
 - ✓ Describe the process of performing non-restoring division

Integer Division

- Unlike subtraction, division cannot easily be performed using multiplication since it is difficult to take the reciprocal of a number.
- Division is performed by repeated subtraction
- There are two methods for integer binary division
 - Restoring
 - Non-restoring

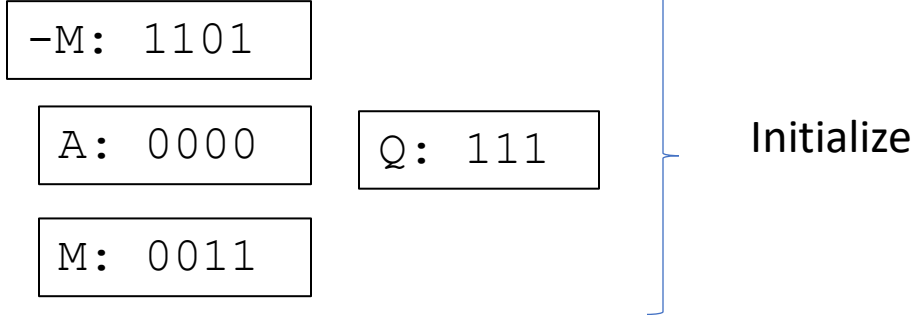
Non-restoring Method

- The non-restoring division defers restoration to reduce operations
 - Addition is performed instead of subtraction in the next step
 - This requires a final restoration, when last operation yielded a negative remainder

Non- Restoring Method

- Initialization
 - Clear A . Requires 1 extra bit for A to be used as a sign bit.
 - Q gets dividend.
 - M gets divisor.
- Loop for each bit of the dividend Q
 - If A (previous result) is negative ($A_n = 1$)
 - Shift AQ to the left.
 - Add ($A \leftarrow A + M$)
 - Q_0 gets the complement value of the MSB of A
 - Else
 - Shift AQ to the left.
 - Subtract ($A \leftarrow A - M$)
 - Q_0 gets the complement value of the MSB of A
- Perform restoration if last result in A is negative, *i.e.*, if $A_n = 1$ then $A \leftarrow A + M$
- Quotient in Q while remainder in A ; adjust sign as needed

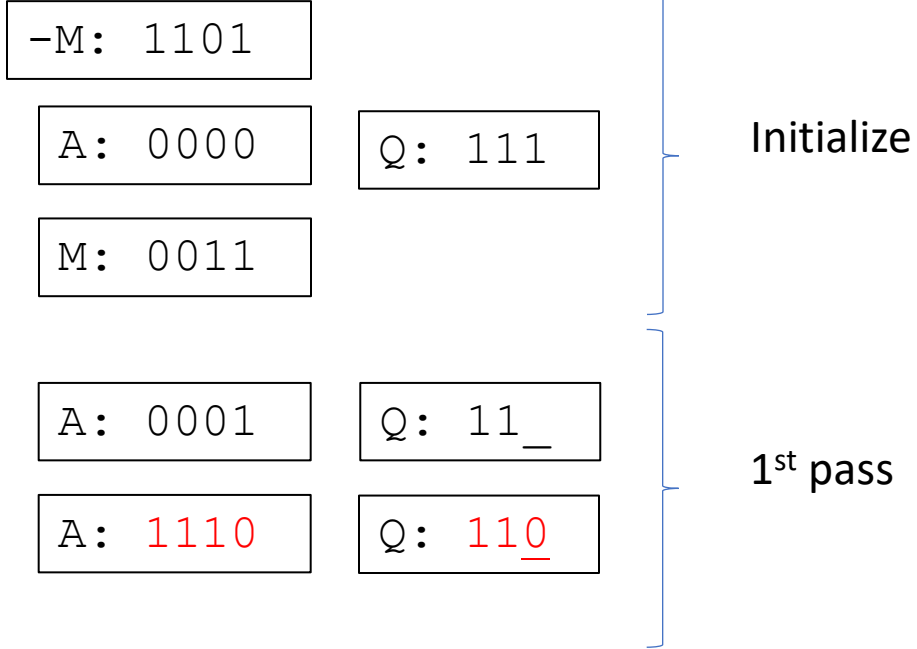
Non-Restoring Division



111 (Q) / 11 (M)

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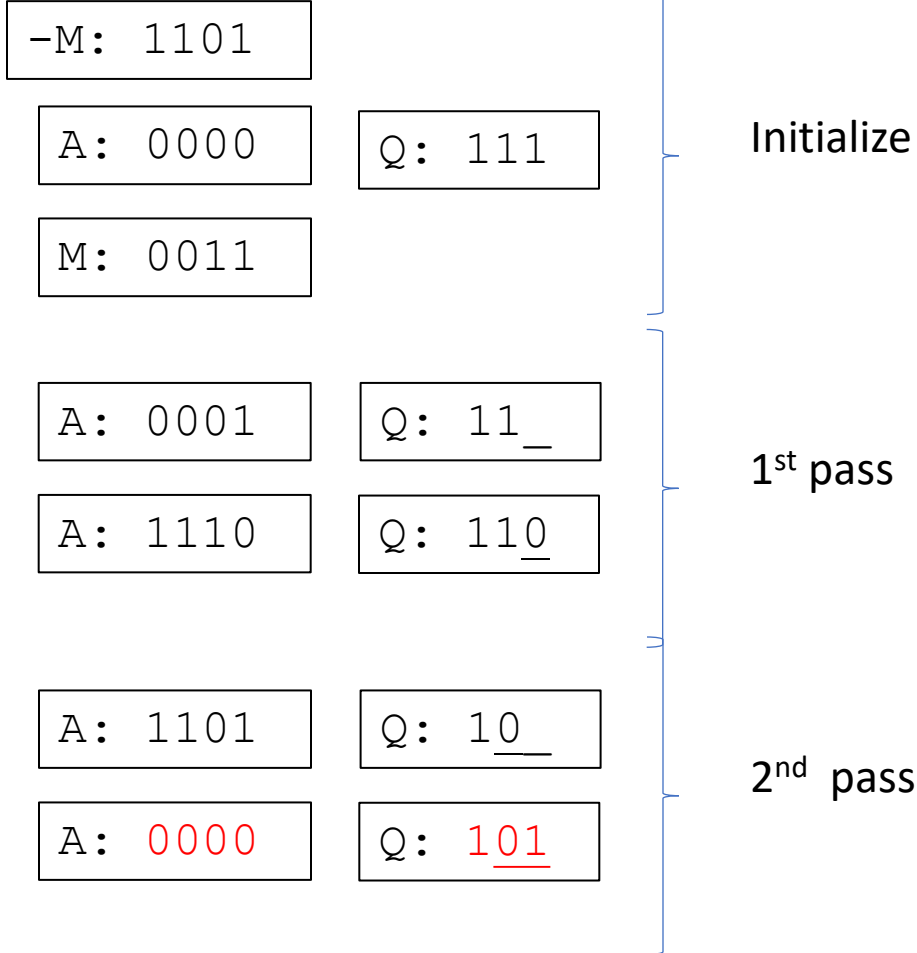
Non-Restoring Division



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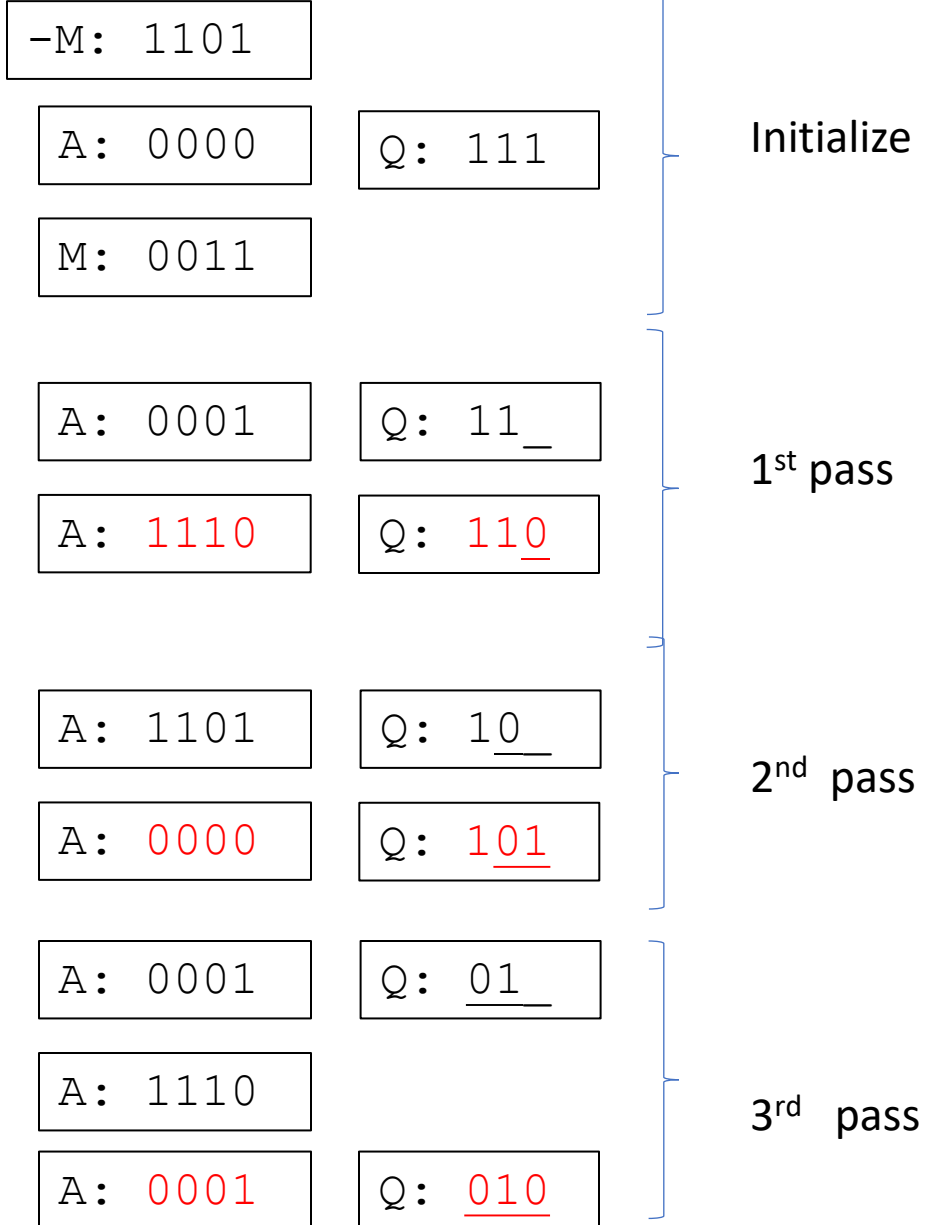
Non-Restoring Division



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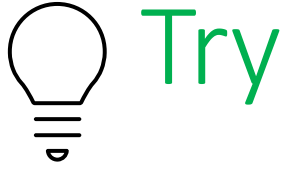
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Non-Restoring Division



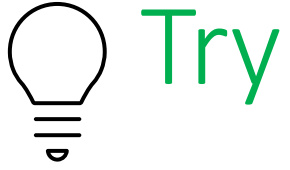
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Try: 01101 (Q) / 00101 (M) (using non-restoring division)
Show the value of A and Q after the end of each pass

After this pass	A	Q
1 st		
2 nd		
3 rd		
4 th		
5 th		



Try: 01101 (Q) / 00101 (M) (using non-restoring division)
Show the value of A and Q after the end of each pass

After this pass	A	Q
1 st	111011	11010
2 nd	111100	10100
3 rd	11110	01000
4 th	000001	10001
5 th	000011	00010

To recall ...

- What have we learned:
 - ✓ Describe the process of performing non-restoring division