A 3-player game theoretic model of a choice between two queueing systems with strategic managerial decision making

Michalis Panayides

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About me

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THIS.

Queues - M/M/1

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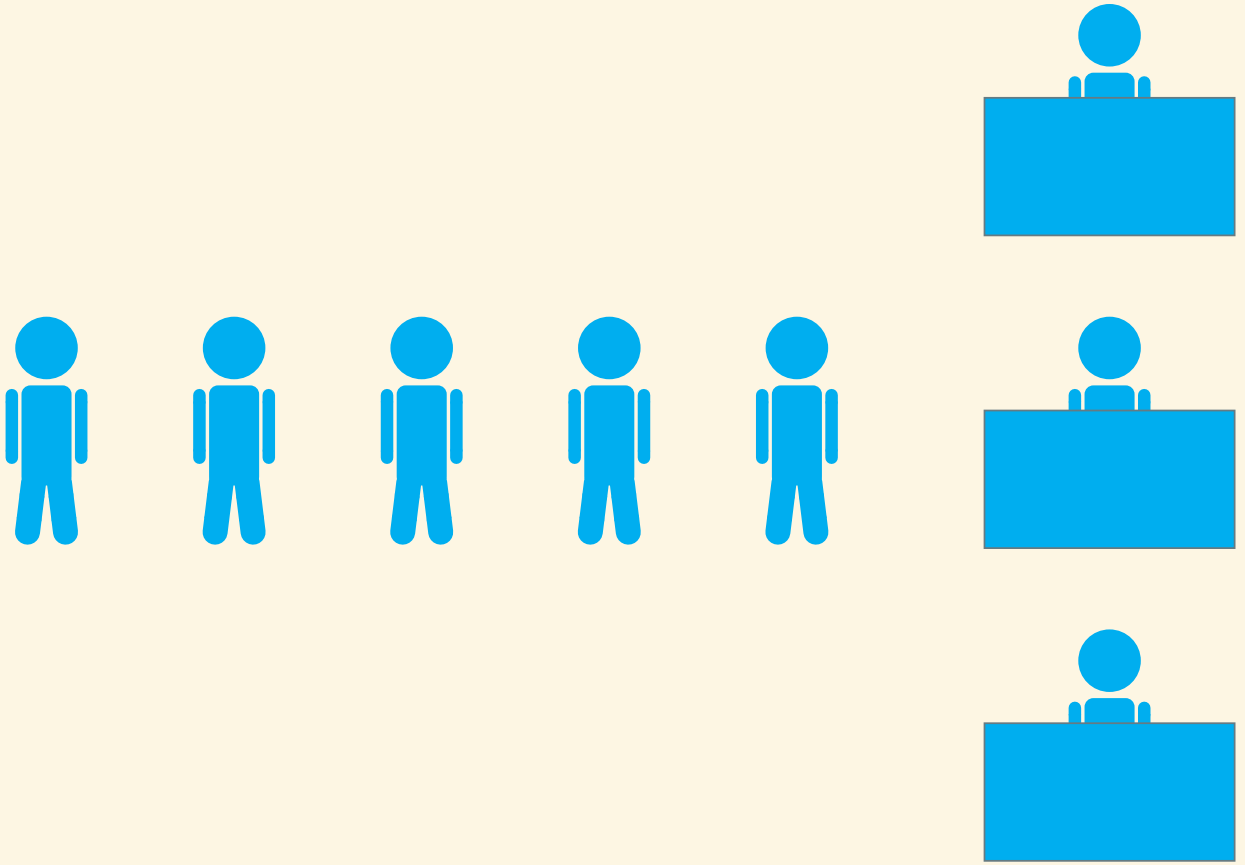
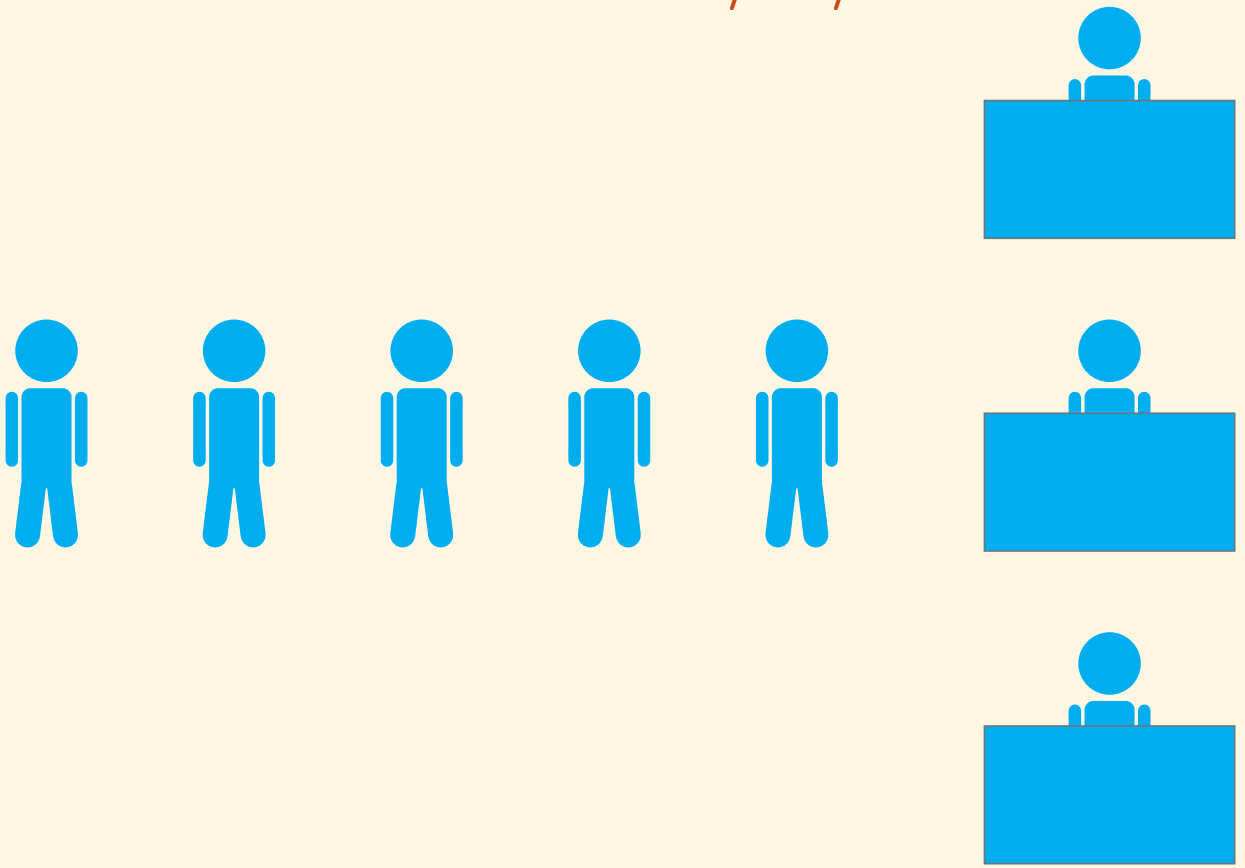
- Shone R, Knight VA, Williams JE. Comparisons between observable and unobservable M/M/1 queues with respect to optimal customer behaviour
- Kerner Y, Shmuel-Bittner O. Strategic behaviour and optimization in a hybrid M/M/1 queue with retrials.
- Gai Y, Liu H, Krishnamachari B. A packet dropping mechanism for efficient operation of M/M/1 queues with selfish users.

Queues - M/M/3



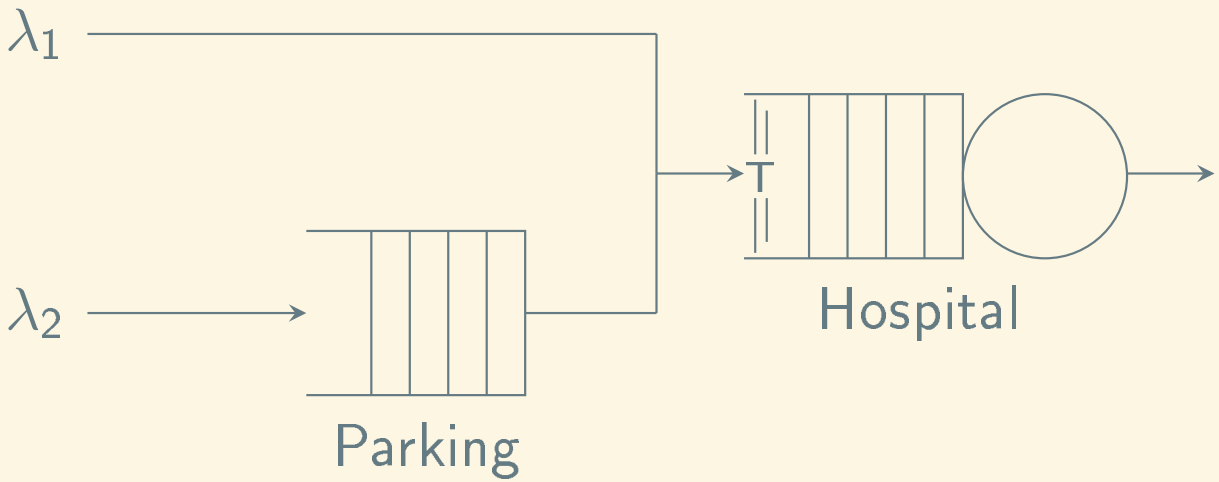
Queues - M/M/3



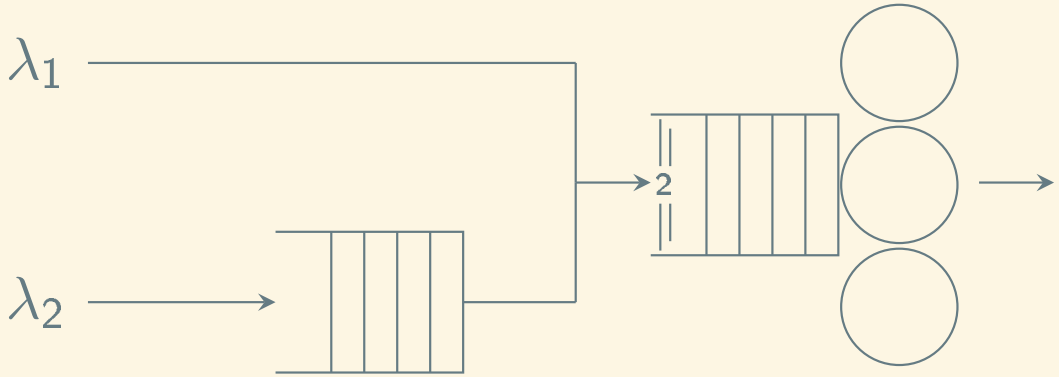
- Knight V, Harper P. The Impact of Choice on Public Services.
- Haviv M, Roughgarden T. The price of anarchy in an exponential multi-server
- Shone R, Knight V, Harper P. A conservative index heuristic for routing problems with multiple heterogeneous service facilities

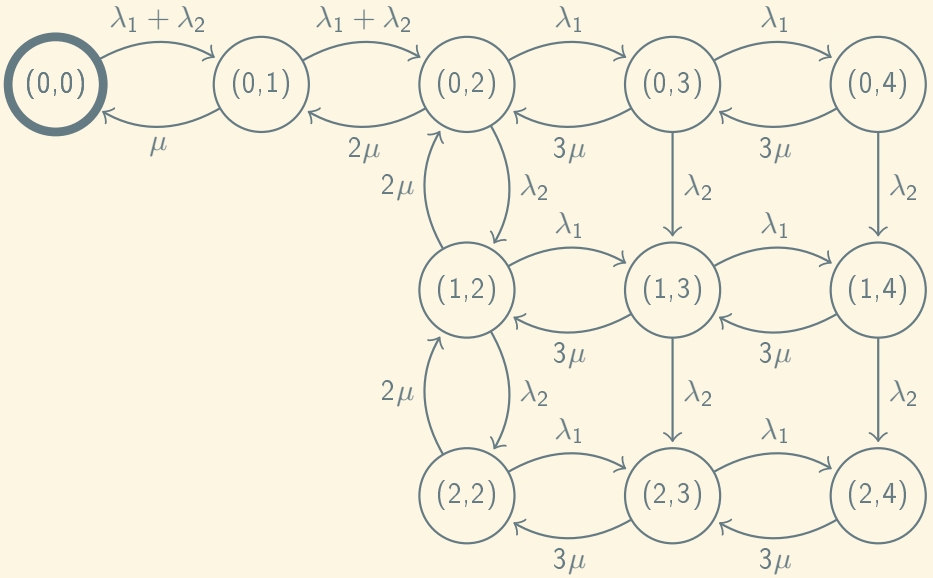
Queues - Custom network of queues

Queueing network structure



Markov Chain - Custom network





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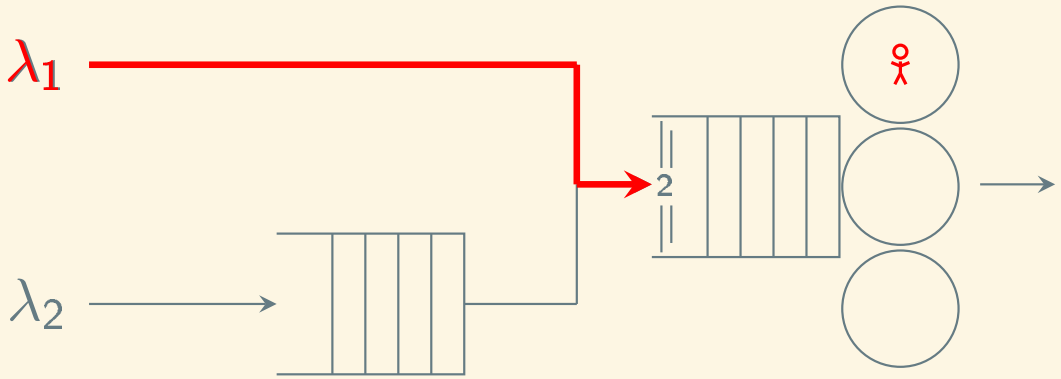

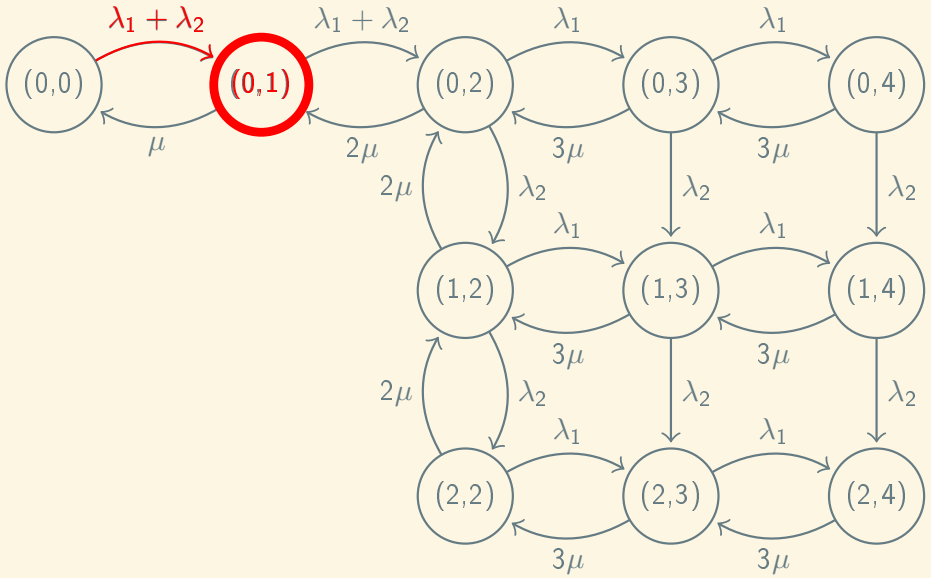

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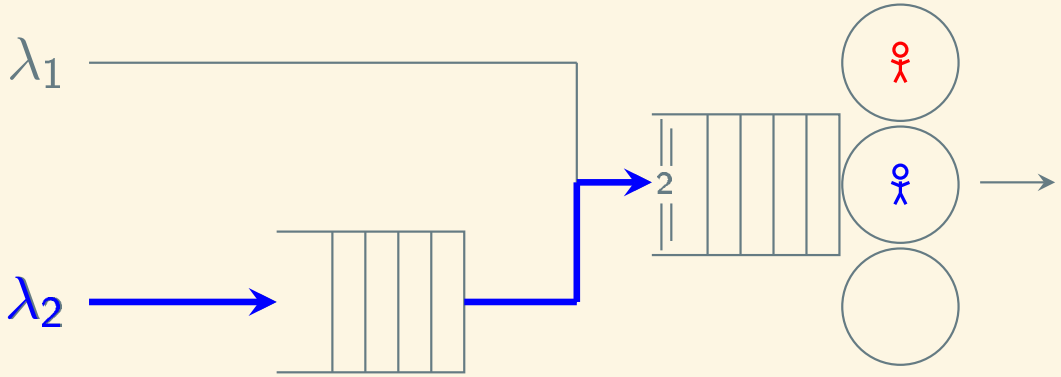

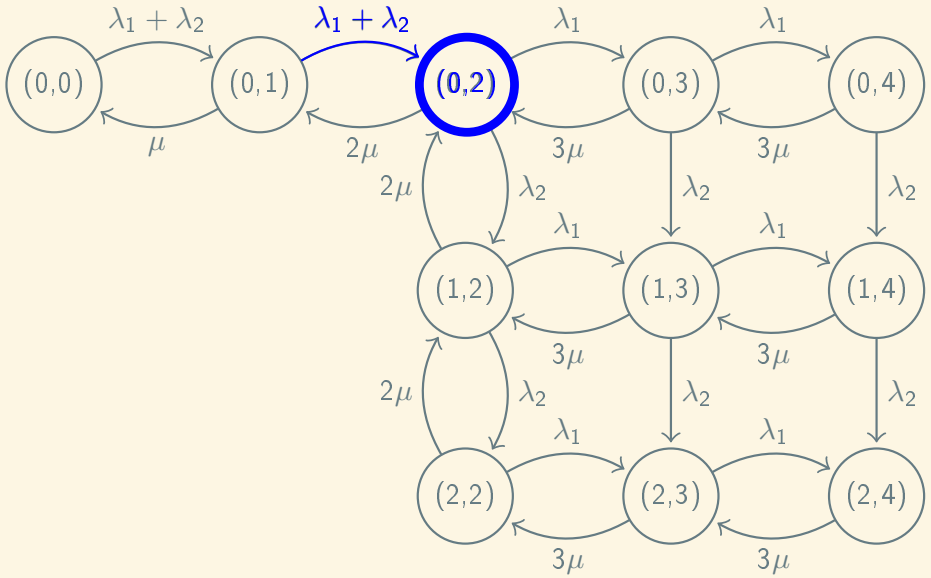

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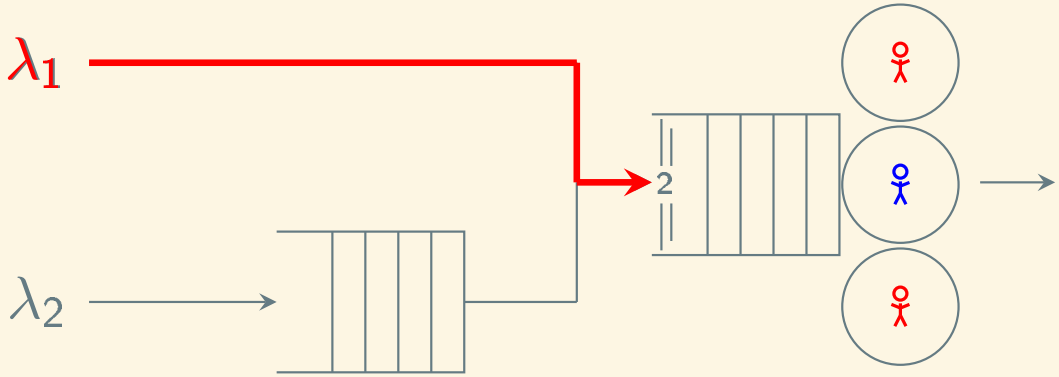

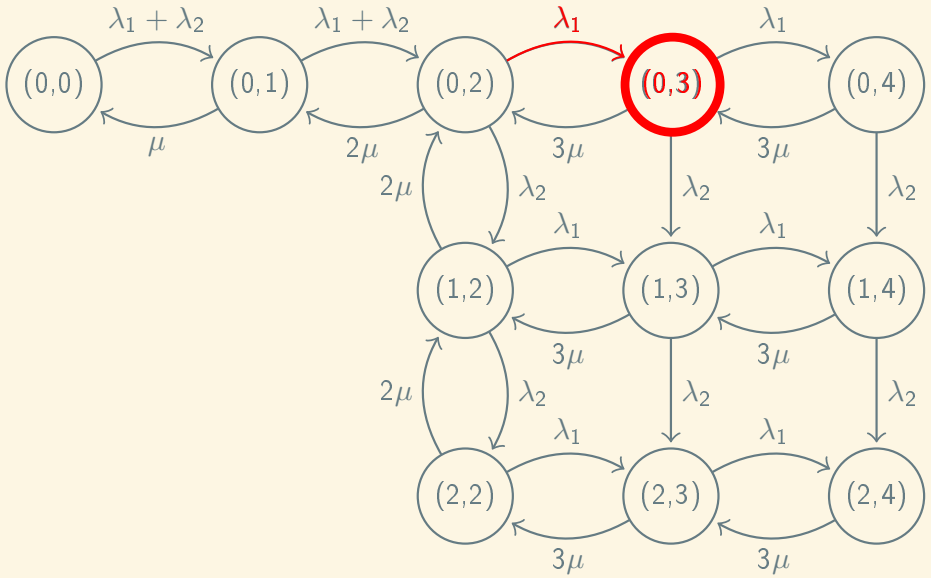

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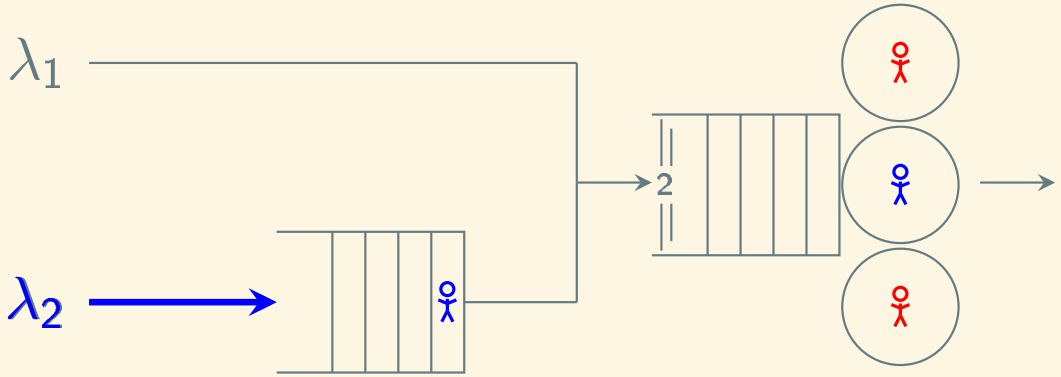

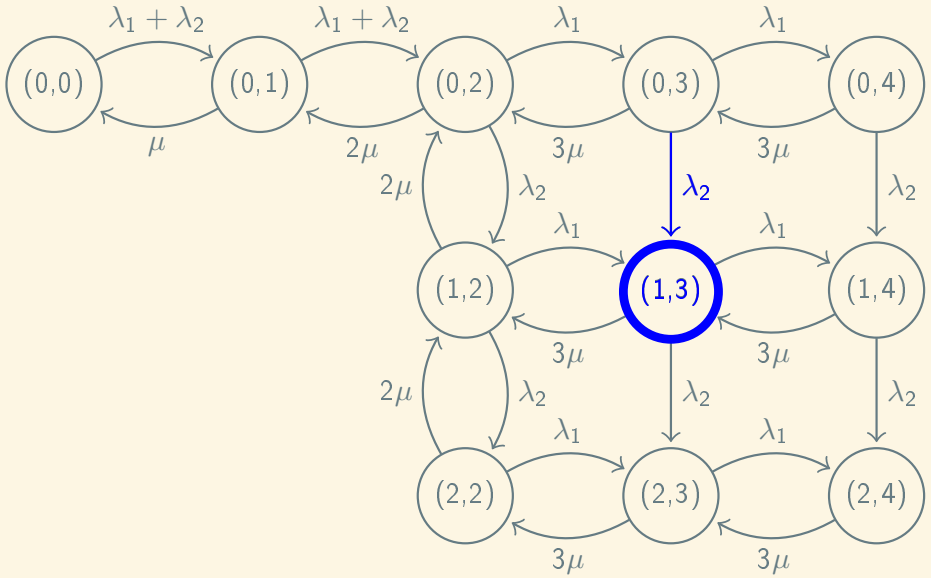

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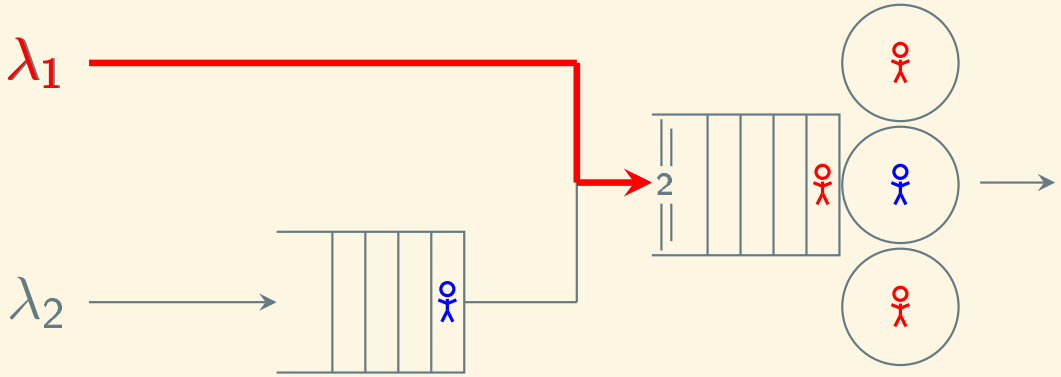
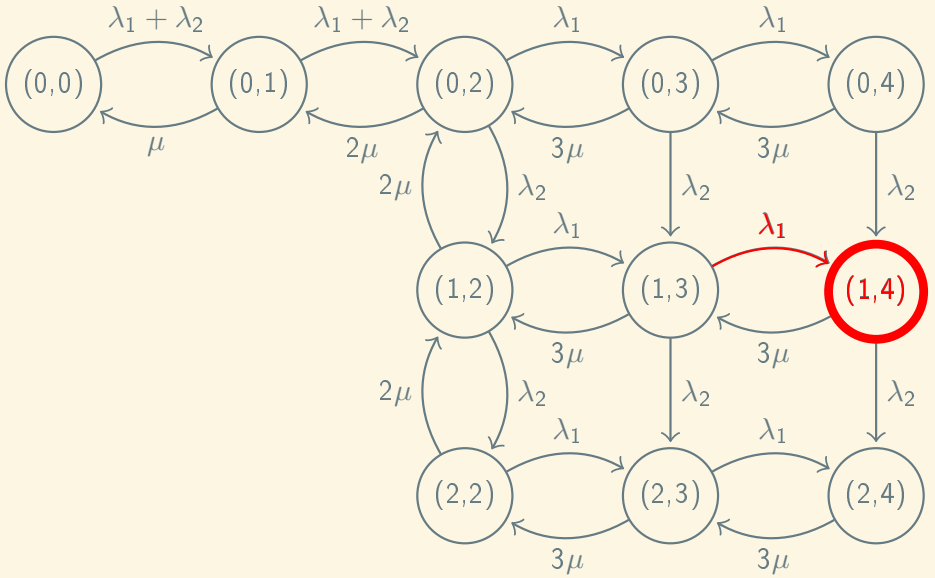
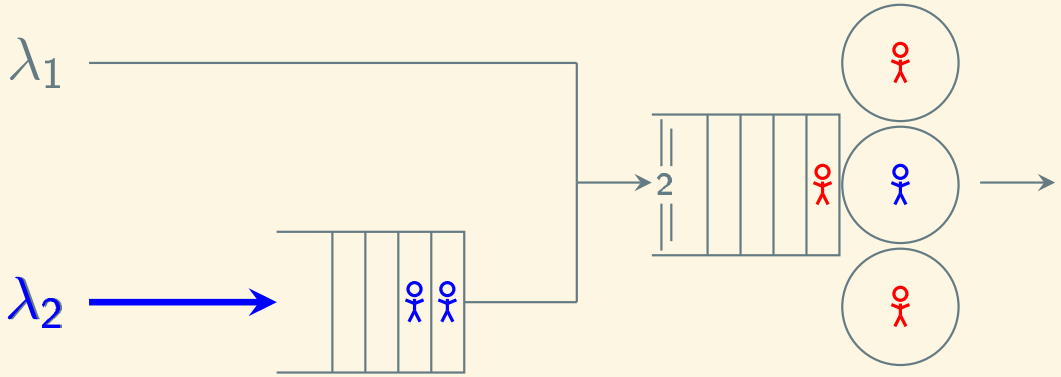
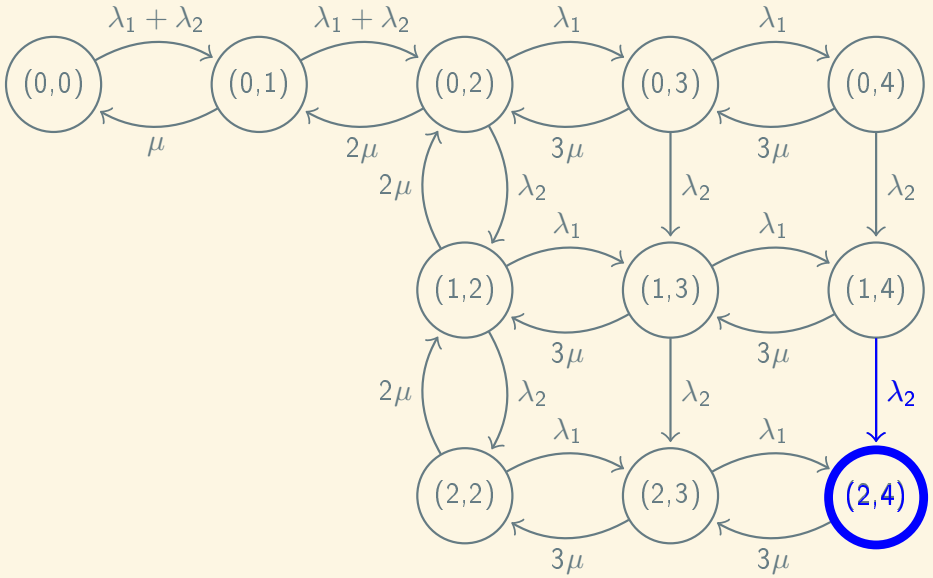




Markov Chain - Custom network





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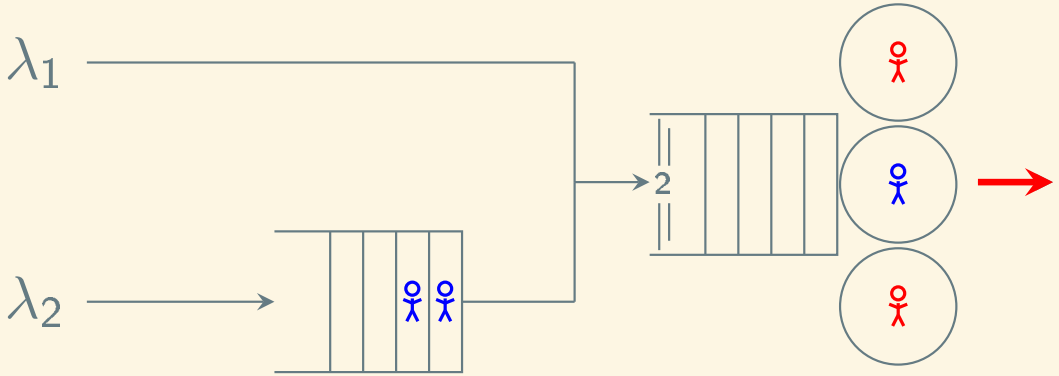

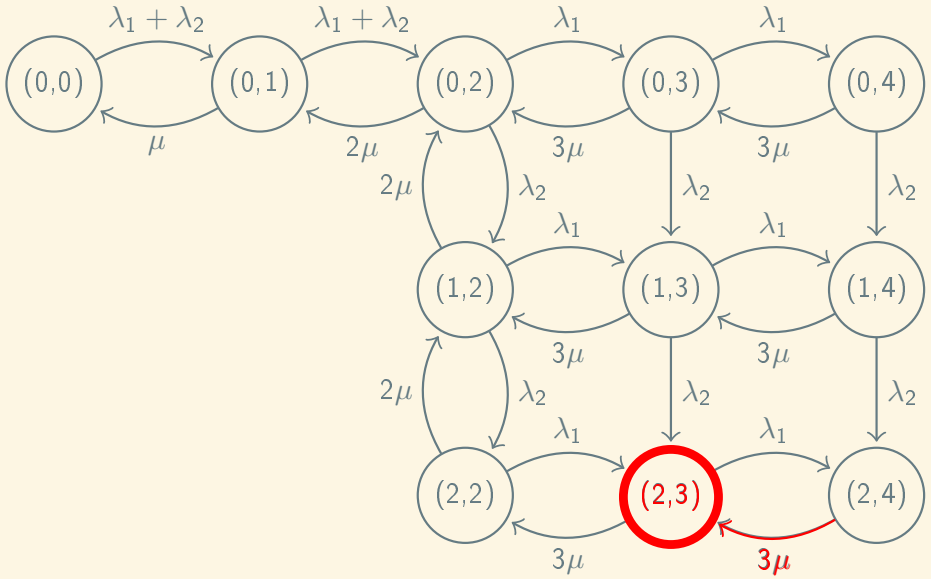

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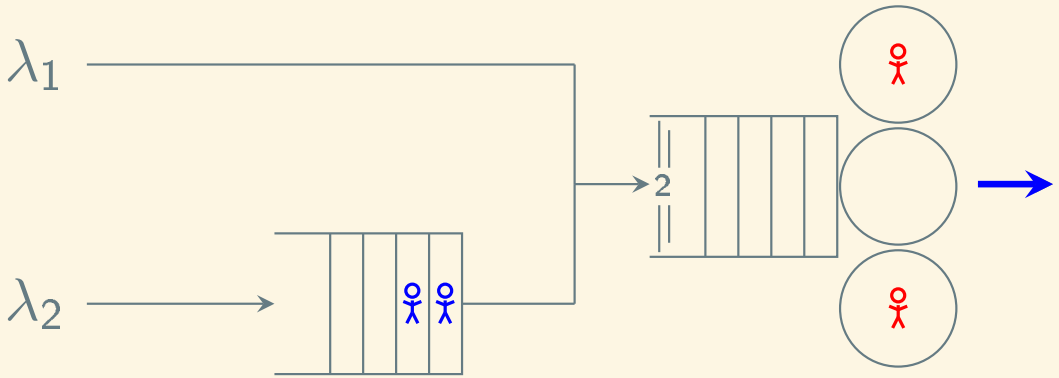
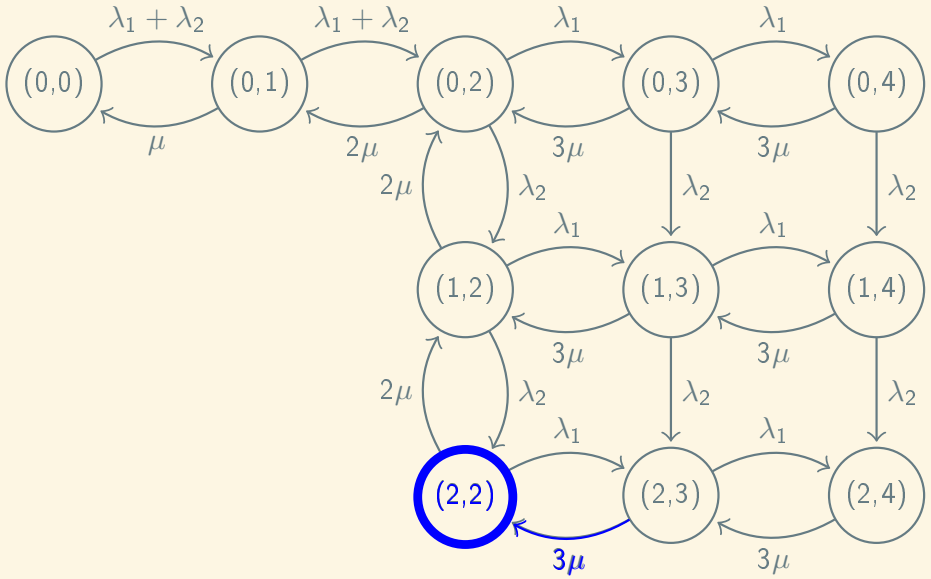
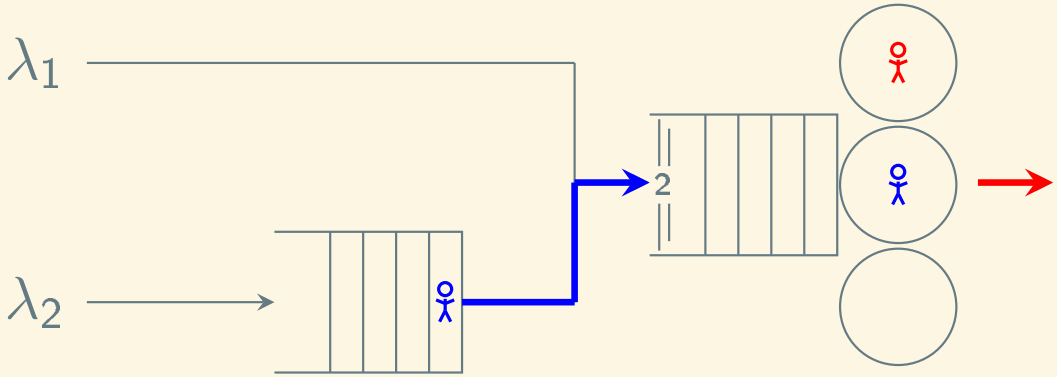
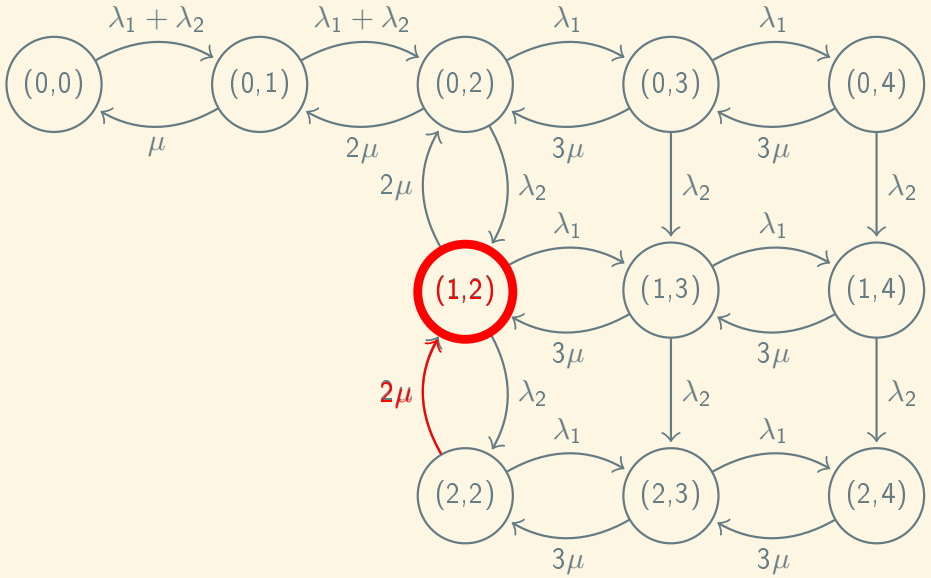




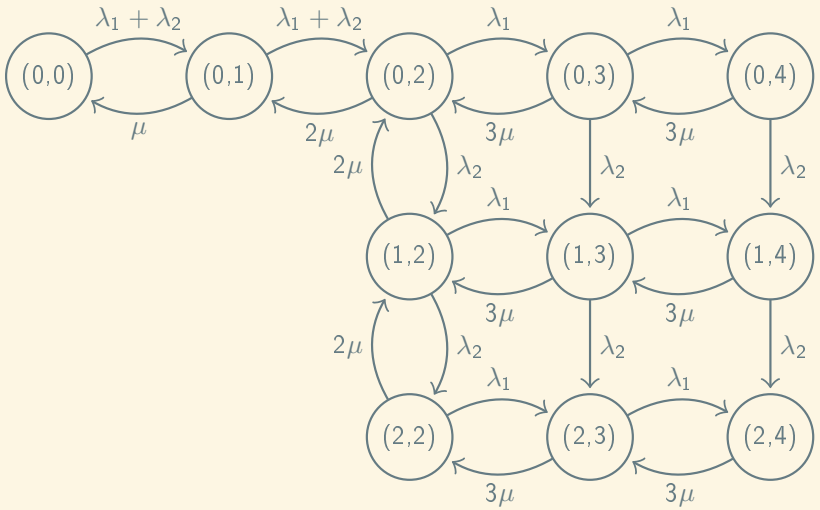
Markov Chain - Custom network





Steady state probabilities - Custom network

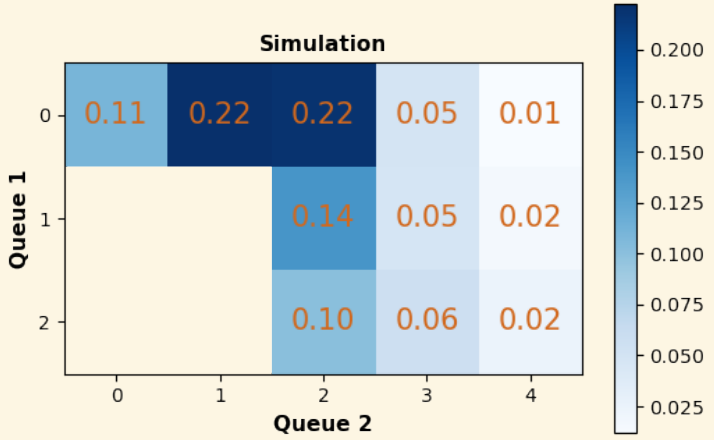
$$
Q = \begin{array}{c} \begin{matrix} (0,0) & (0,1) & (0,2) & \dots & (2,3) & (2,4) \end{matrix} \\ \begin{pmatrix} -\lambda_1 - \lambda_2 & \lambda_1 + \lambda_2 & 0 & \dots & 0 & 0 \\ \mu & -\mu - \lambda_1 - \lambda_2 & \lambda_1 + \lambda_2 & \dots & 0 & 0 \\ 0 & 2\mu & -2\mu - \lambda_1 - \lambda_2 & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & -\lambda_1 - 3\mu & \lambda_1 \\ 0 & 0 & 0 & \dots & 3\mu & -3\mu \end{pmatrix} \end{array} \begin{matrix} (0,0) \\ (0,1) \\ (0,2) \\ \\ \\ (2,3) \\ (2,4) \end{matrix}
$$

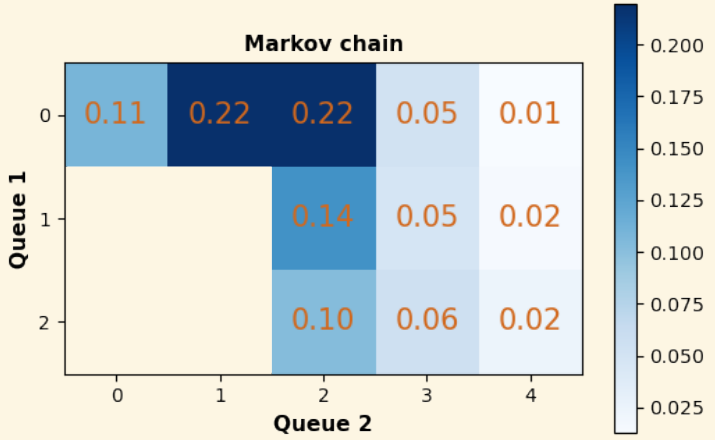


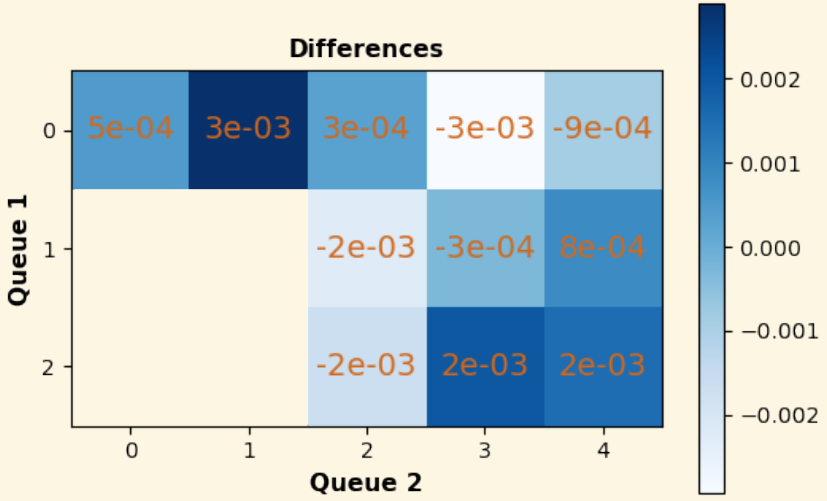
$$
\frac{d\pi}{dt} = \pi Q = 0, \qquad \sum \pi_{(u,v)} = 1
$$

$$
\pi = \begin{bmatrix} \pi_{(0,0)} \\ \pi_{(0,1)} \\ \pi_{(0,2)} \\ \vdots \\ \pi_{(2,3)} \\ \pi_{(2,4)} \end{bmatrix}
$$

Steady state probabilities - Comparison





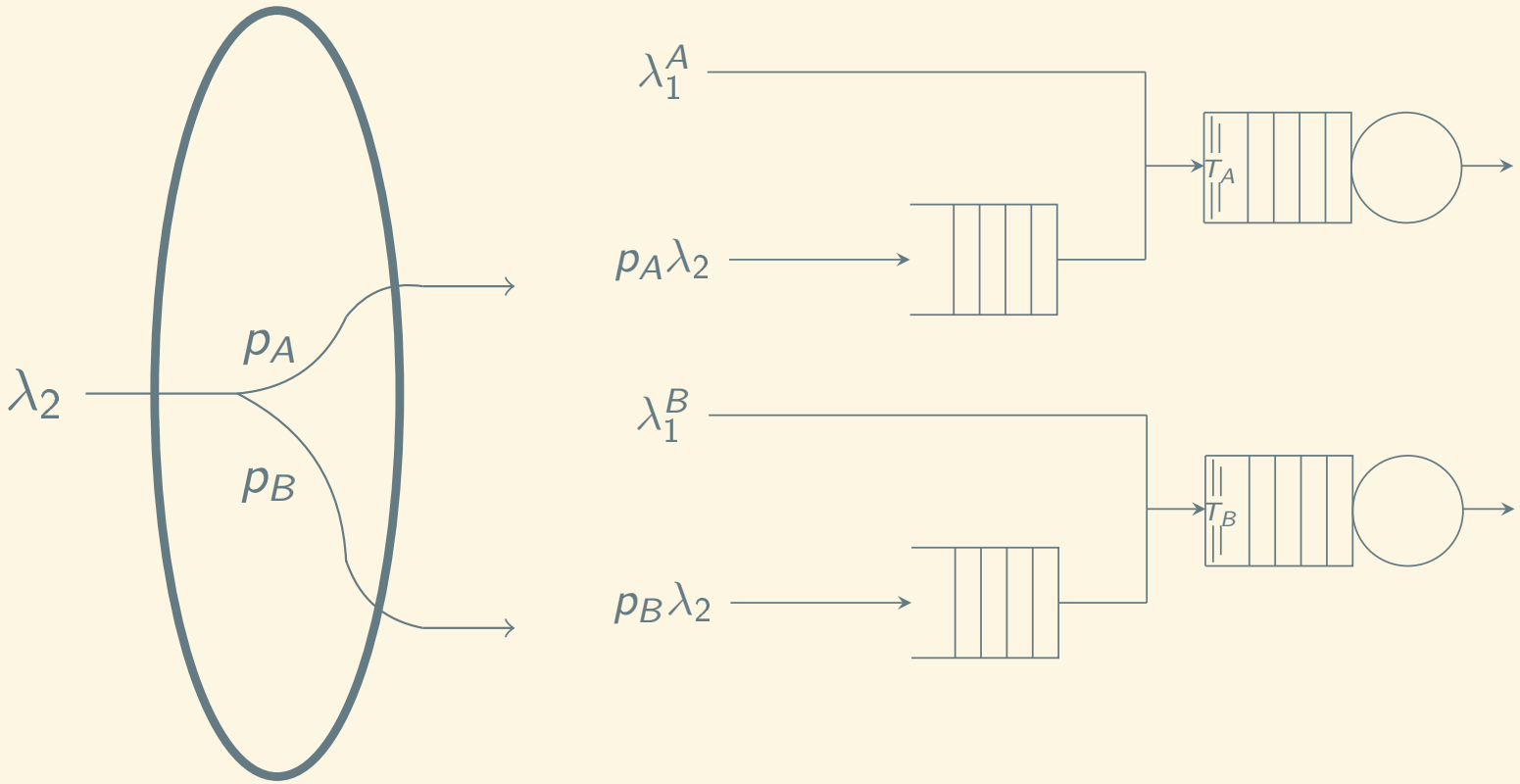


Game - Definition

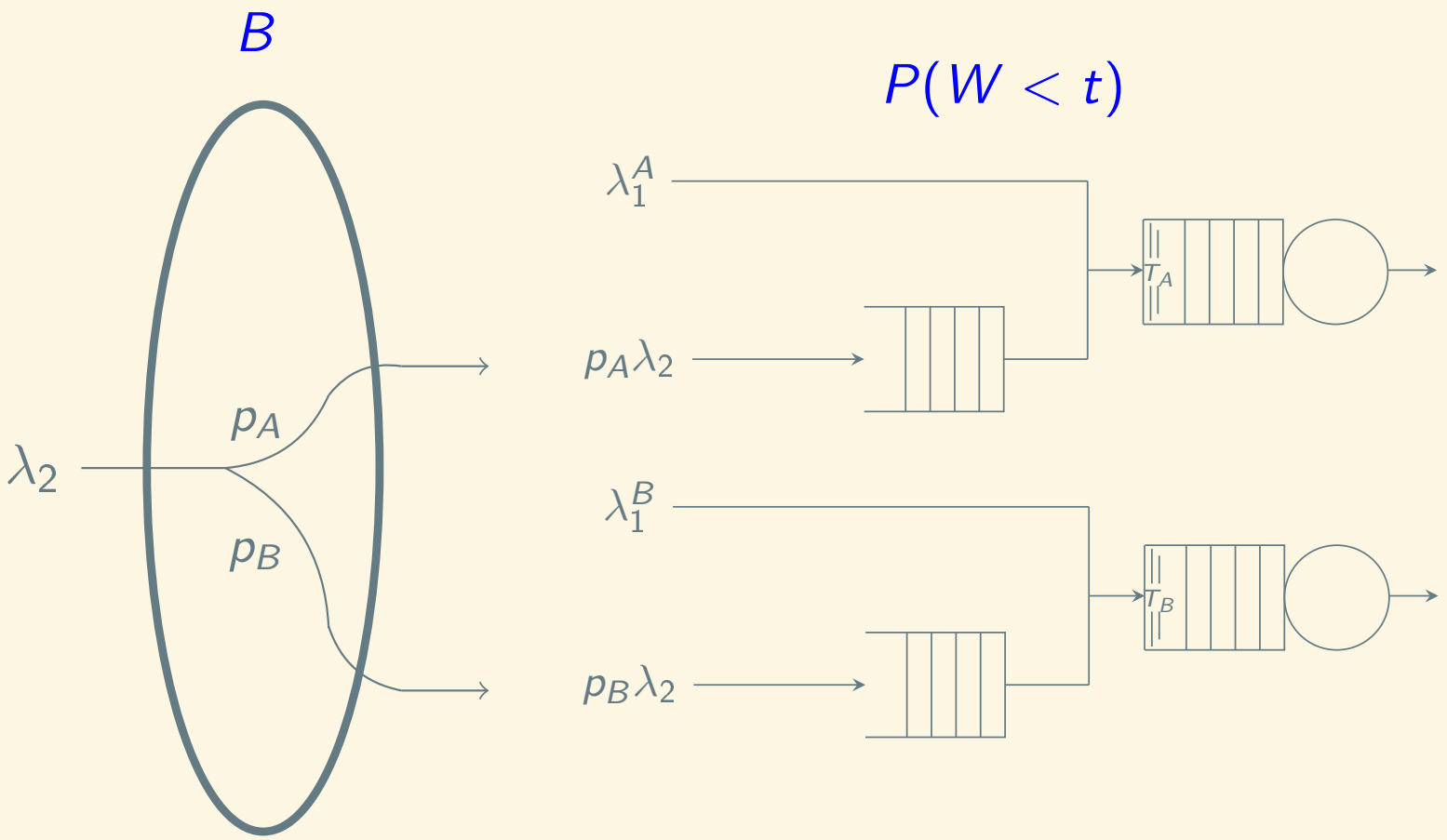
Game - Definition

- Misra S, Sarkar S. Priority-based time-slot allocation in wireless body area networks during medical emergency situations: An evolutionary game-theoretic perspective
- Song J, Wen J. A non-cooperative game with incomplete information to improve patient hospital choice

Game - Players and objectives



Game - Players and objectives



Performance Measures - Blocking time

$$B = \frac{\sum_{(u,v) \in S_A^{(2)}} \pi_{(u,v)} \; b(u,v)}{\sum_{(u,v) \in S_A^{(2)}} \pi_{(u,v)}}$$

Performance Measures - Blocking time

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Performance Measures - Blocking time

$$B = \frac{\sum_{(u,v)\in S_A^{(2)}} \pi_{(u,v)}\ b(u,v)}{\sum_{(u,v)\in S_A^{(2)}} \pi_{(u,v)}}$$

$$b(u,v) = \begin{cases} 0, & \text{if } (u,v) \notin S_b \\ c(u,v) + b(u-1,v), & \text{if } v = N = T \\ c(u,v) + b(u,v-1), & \text{if } v = N \neq T \\ c(u,v) + p_s(u,v)b(u-1,v) + p_a(u,v)b(u,v+1), & \text{if } u > 0 \text{ and } v = T \\ c(u,v) + p_s(u,v)b(u,v-1) + p_a(u,v)b(u,v+1), & \text{otherwise} \end{cases}$$

$$S_b = \{(u,v) \in S \mid u > 0\}$$

$$c(u,v) = \begin{cases} \frac{1}{\min(v,C)\mu}, & \text{if } v = N \\ \frac{1}{\lambda_1 + \min(v,C)\mu}, & \text{otherwise} \end{cases}$$

$$p_s(u,v) = \frac{\min(v,C)\mu}{\lambda_1 + \min(v,C)\mu}, \qquad p_a(u,v) = \frac{\lambda_1}{\lambda_1 + \min(v,C)\mu}$$

Performance Measures - Proportion within time

$$P(W < t) = \frac{\lambda_1 P_{L_1'}}{\lambda_2 P_{L_2'} + \lambda_1 P_{L_1'}} P(W^{(1)} < t) + \frac{\lambda_2 P_{L_2'}}{\lambda_2 P_{L_2'} + \lambda_1 P_{L_1'}} P(W^{(2)} < t)$$

$$P(W^{(1)} < t) = \frac{\sum_{(u,v) \in S_A^{(1)}} P(W_{(u,v)}^{(1)} < t) \pi_{u,v}}{\sum_{(u,v) \in S_A^{(1)}} \pi_{u,v}}$$

$$P(W^{(2)} < t) = \frac{\sum_{(u,v) \in S_A^{(2)}} P(W_{(u,v)}^{(2)} < t) \pi_{u,v}}{\sum_{(u,v) \in S_A^{(2)}} \pi_{u,v}}$$

Performance Measures - Proportion within time

$$P(W^{(i)} < t) = \frac{\sum_{(u,v)\in S_A^{(i)}} P(W_{u,v}^{(i)} < t)\pi_{u,v}}{\sum_{(u,v)\in S_A^{(i)}} \pi_{u,v}}, \quad \textbf{for } i = \{1, 2\}$$

Performance Measures - Proportion within time

$$P(W^{(i)} < t) = \frac{\sum_{(u,v)\in S_A^{(i)}} P(W_{u,v}^{(i)} < t)\pi_{u,v}}{\sum_{(u,v)\in S_A^{(i)}} \pi_{u,v}}, \quad \textbf{for } i = \{1, 2\}$$

$$W_{(u,v)}^{(1)} \sim \begin{cases} \textbf{Erlang}(v, \mu), & \textbf{if } C = 1 \textbf{ and } v > 1 \\ \textbf{Hypo}\left([v - C, 1], [C\mu, \mu]\right), & \textbf{if } C > 1 \textbf{ and } v > C \\ \textbf{Erlang}(1, \mu), & \textbf{if } v \leq C \end{cases}$$

$$W_{(u,v)}^{(2)} \sim \begin{cases} \textbf{Erlang}(\min(v, T), \mu), & \textbf{if } C = 1 \textbf{ and } v, T > 1 \\ \textbf{Hypo}\left([\min(v, T) - C, 1], [C\mu, \mu]\right), & \textbf{if } C > 1 \textbf{ and } v, T > C \\ \textbf{Erlang}(1, \mu), & \textbf{if } v \leq C \textbf{ or } T \leq C \end{cases}$$

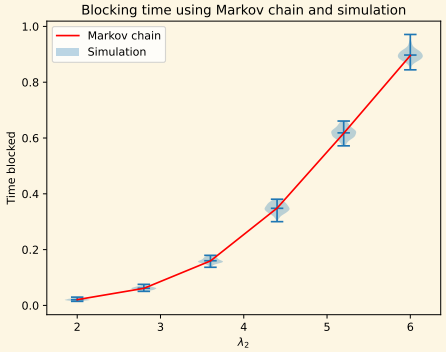
Performance Measures - Proportion within time

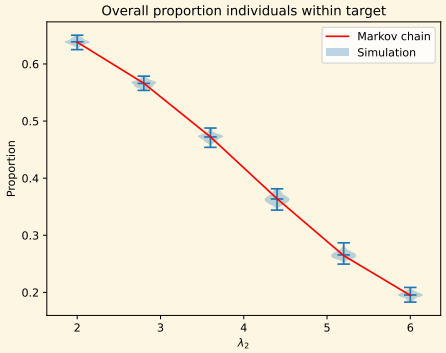
$$P(W^{(i)} < t) = \frac{\sum_{(u,v)\in S_A^{(i)}} P(W_{u,v}^{(i)} < t)\pi_{u,v}}{\sum_{(u,v)\in S_A^{(i)}} \pi_{u,v}}, \quad \textbf{for } i = \{1, 2\}$$

$$P(W_{(u,v)}^{(\mathbf{1})} < t) = \begin{cases} 1 - \sum_{i=\mathbf{0}}^{v-\mathbf{1}} \frac{\mathbf{1}}{i!} e^{-\mu t}(\mu t)^i, & \text{if } C = 1 \text{ and } v > 1 \\ 1 - (\mu C)^{v-C} \mu \sum_{k=\mathbf{1}}^{|\vec{r}|} \sum_{l=\mathbf{1}}^{r_k} \frac{\Psi_{k,l}(-\lambda_k) t^{r_k - l} e^{-\lambda_k t}}{(r_k - l)!(l-\mathbf{1})!}, & \text{if } C > 1 \text{ and } v > C \\ \quad \textbf{where } \vec{r} = (v - C, 1) \textbf{ and } \vec{\lambda} = (C\mu, \mu) & \\ 1 - e^{-\mu t}, & \text{if } v \leq C \end{cases}$$

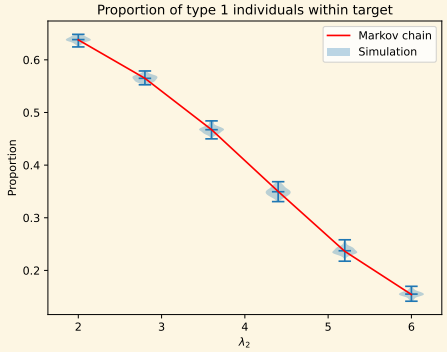
$$P(W_{(u,v)}^{(\mathbf{2})} < t) = \begin{cases} 1 - \sum_{i=\mathbf{0}}^{\min(v,T)-\mathbf{1}} \frac{\mathbf{1}}{i!} e^{-\mu t}(\mu t)^i, & \text{if } C = 1 \text{ and } v, T > 1 \\ 1 - \mu(\mu C)^{\min(v,T)-C} \sum_{k=\mathbf{1}}^{|\vec{r}|} \sum_{l=\mathbf{1}}^{r_k} \frac{\Psi_{k,l}(-\lambda_k) t^{r_k - l} e^{-\lambda_k t}}{(r_k - l)!(l-\mathbf{1})!}, & \text{if } C > 1 \text{ and } v, T > C \\ \quad \textbf{where } \vec{r} = (\min(v, T) - C, 1) \textbf{ and } \vec{\lambda} = (C\mu, \mu) & \\ 1 - e^{-\mu t}, & \text{if } v \leq C \textbf{ or } T \leq C \end{cases}$$

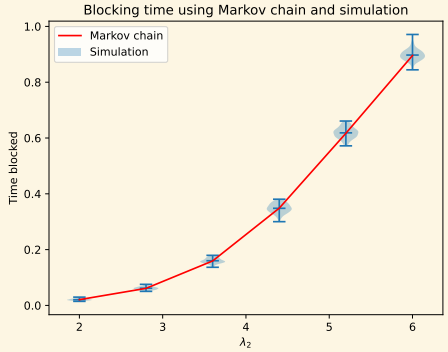
Comparisons

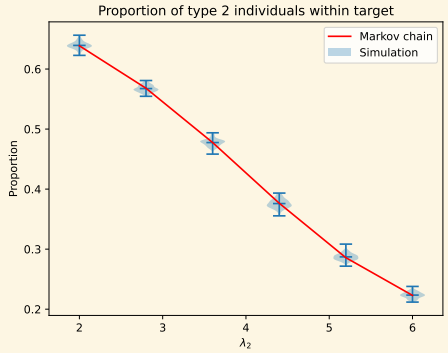




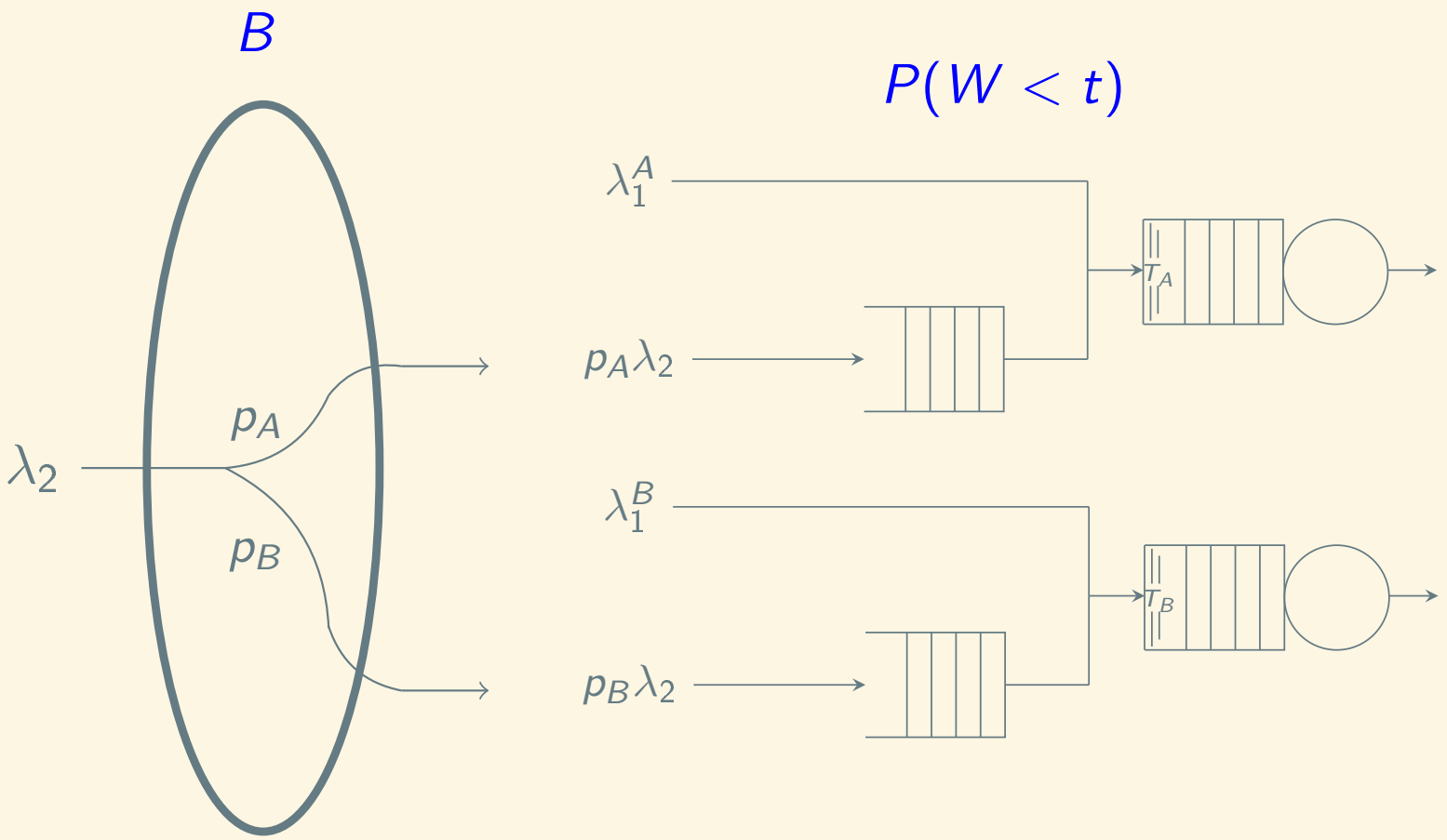
N/A



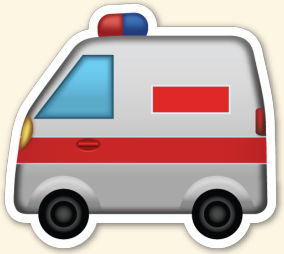


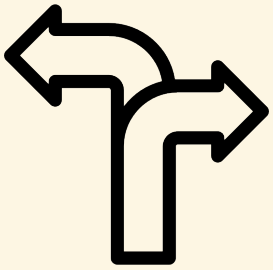


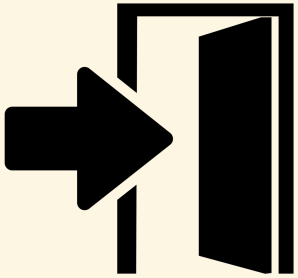
Game - Players and objectives

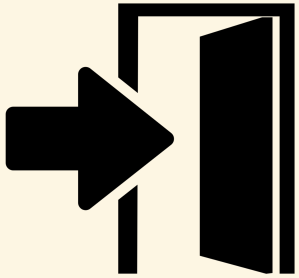


Game - Strategies









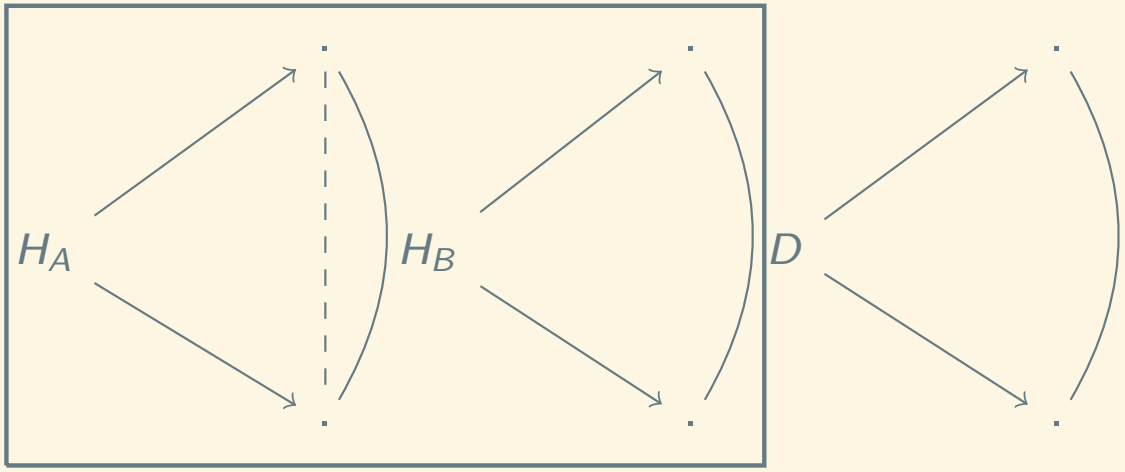
$p_A, p_B \in [0, 1]$

$T_A \in [1, N_A]$

$T_B \in [1, N_B]$

$p_A + p_B = 1$

Game - Formulation

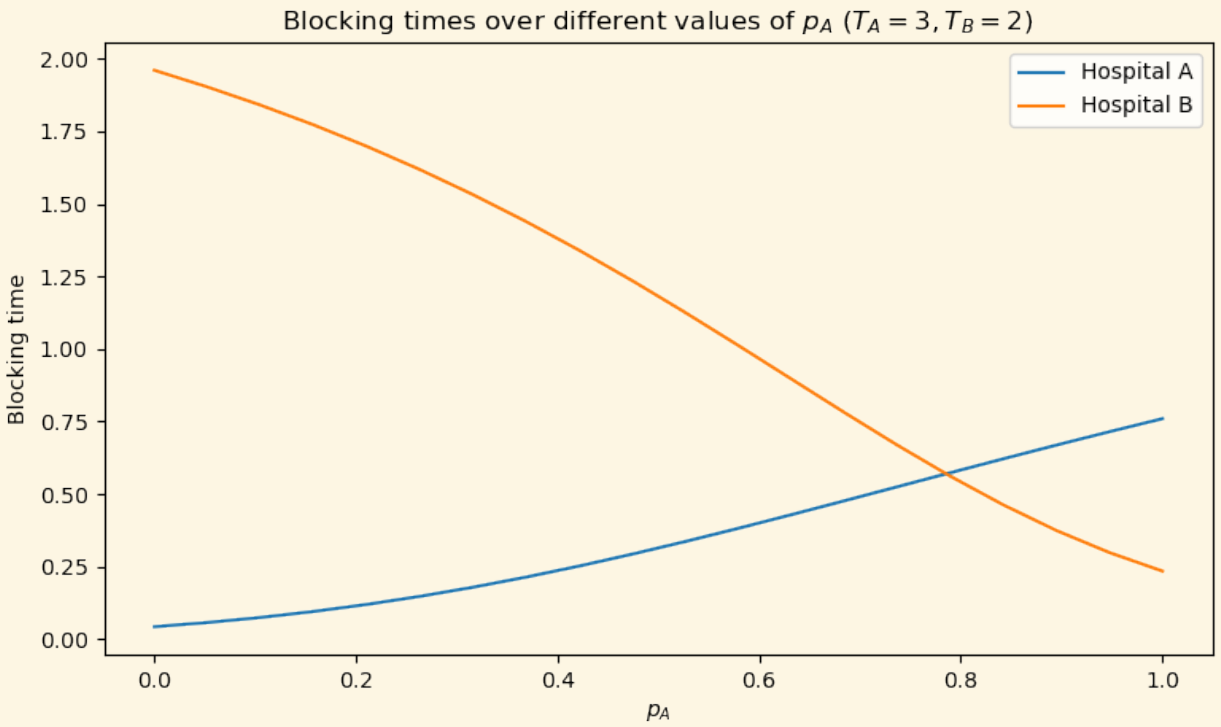


Ambulance's Decision

$$R = \begin{pmatrix} p_{1,1} & p_{1,2} & \cdots & p_{1,N_B} \\ p_{2,1} & p_{2,2} & \cdots & p_{2,N_B} \\ \vdots & \vdots & \ddots & \vdots \\ p_{N_A,1} & p_{N_A,2} & \cdots & p_{N_A,N_B} \end{pmatrix}$$

Ambulance's Decision

$$\begin{matrix} T_A = 3 \\ T_B = 2 \end{matrix} \rightarrow \begin{pmatrix} - & - & - & - \\ - & - & - & - \\ - & x & - & - \\ - & - & - & - \end{pmatrix} \rightarrow B_A(p_A) = B_B(1 - p_A)$$



Hospitals' Decision

$$A = \begin{pmatrix} U^A_{1,1} & U^A_{1,2} & \cdots & U^A_{1,N_B} \\ U^A_{2,1} & U^A_{2,2} & \cdots & U^A_{2,N_B} \\ \vdots & \vdots & \ddots & \vdots \\ U^A_{N_A,1} & U^A_{N_A,2} & \cdots & U^A_{N_A,N_B} \end{pmatrix}, \quad B = \begin{pmatrix} U^B_{1,1} & U^B_{1,2} & \cdots & U^B_{1,N_B} \\ U^B_{2,1} & U^B_{2,2} & \cdots & U^B_{2,N_B} \\ \vdots & \vdots & \ddots & \vdots \\ U^B_{N_A,1} & U^B_{N_A,2} & \cdots & U^B_{N_A,N_B} \end{pmatrix}$$

$$U^{(i)}_{T_A,T_B} = 1 - \left[(P(X^{(i)} < t) - 0.95)^2 \right]$$

Nash Equilibrium

$$R = \begin{pmatrix} 0.5 & 0.1 & 0 & 0 \\ 0.9 & 0.5 & 0.2 & 0 \\ 1 & 0.8 & 0.5 & 0.3 \\ 1 & 1 & 0.7 & 0.5 \end{pmatrix}$$

$$A = \begin{pmatrix} 0.99998394 & 0.99998394 & 0.99998394 & 0.99998394 \\ 0.99998955 & 0.99998848 & 0.99998649 & 0.9999845 \\ 0.99999952 & 0.9999987 & 0.99999596 & 0.99999199 \\ 0.99994372 & 0.99995113 & 0.99998603 & 0.99999911 \end{pmatrix}$$

$$B = \begin{pmatrix} 0.99998394 & 0.99998955 & 0.99999952 & 0.99994372 \\ 0.99998394 & 0.99998848 & 0.9999987 & 0.99995113 \\ 0.99998394 & 0.99998649 & 0.99999596 & 0.99998603 \\ 0.99998394 & 0.9999845 & 0.99999199 & 0.99999911 \end{pmatrix}$$

Nash Equilibria: $\underline{A}$ $(0, 0, 0.4, 0.6)$ $\underline{B}$ $(0, 0, 0.4, 0.6)$

Asymmetric Replicator Dynamics

$$\frac{dx}{dt}_i = x_i((f_x)_i - \phi_x), \quad \text{for all } i$$

$$\frac{dy}{dt}_i = y_i((f_y)_i - \phi_y), \quad \text{for all } i$$

- Fudenberg, Drew, et al. The theory of learning in games. Vol. 2. MIT press, 1998.
- Elvio, Accinelli and Carrera, Edgar. 2011. Evolutionarily Stable Strategies and Replicator Dynamics in Asymmetric Two-Population Games. 10.1007/978-3-642-11456-4_3.

Inefficiency measure

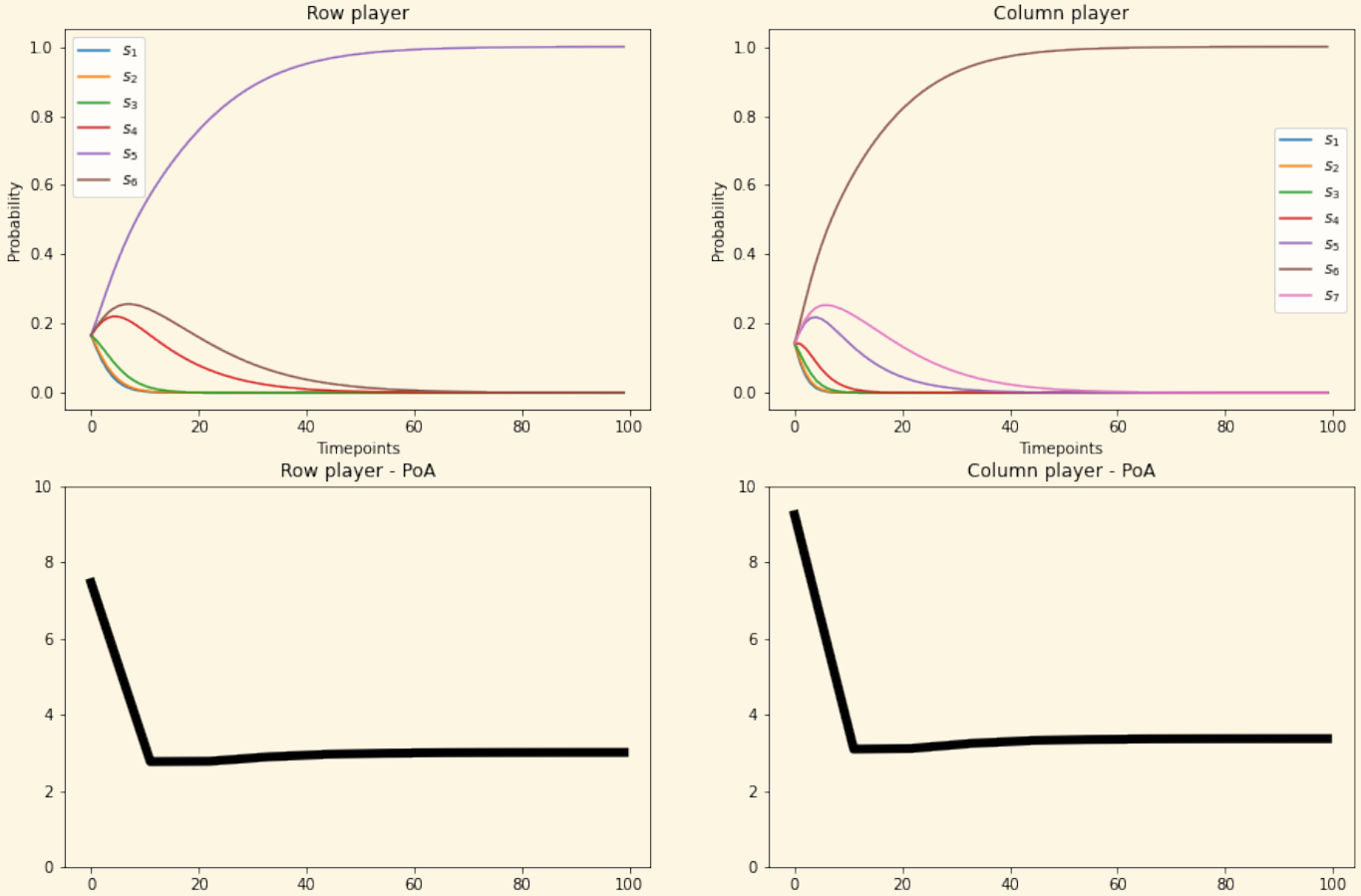
$$PoA = \frac{\max_{s \in E} Cost(s)}{\min_{s \in S} Cost(S)}$$

Inefficiency measure

$$PoA = \frac{\max_{s \in E} Cost(s)}{\min_{s \in S} Cost(S)}$$

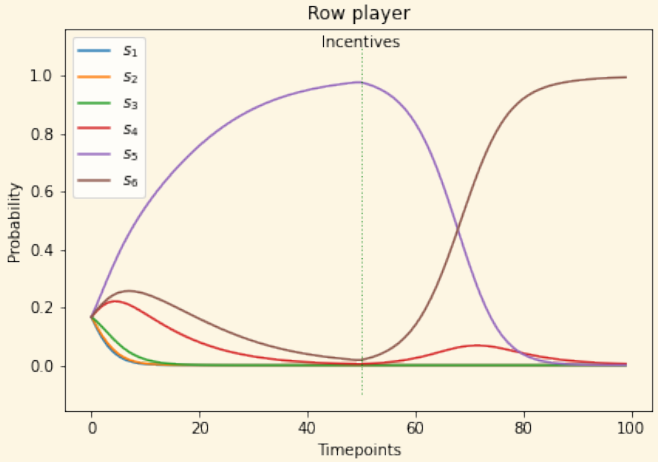
$$PoA_A(s_r) = \frac{Cost(s_r)}{\min_{s \in S} Cost(S)}, \qquad PoA_B(s_c) = \frac{Cost(s_c)}{\min_{s \in S} Cost(S)}$$

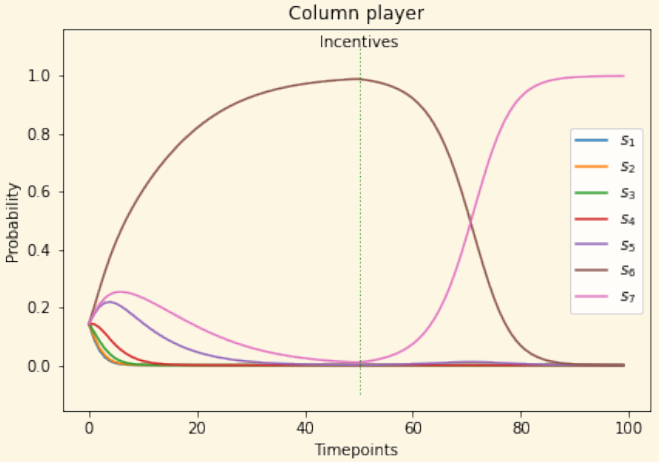
Learning algorithms - Asymmetric replicator dynamics

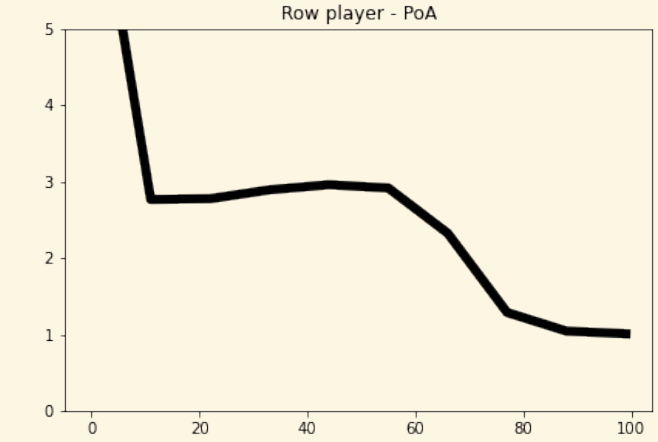


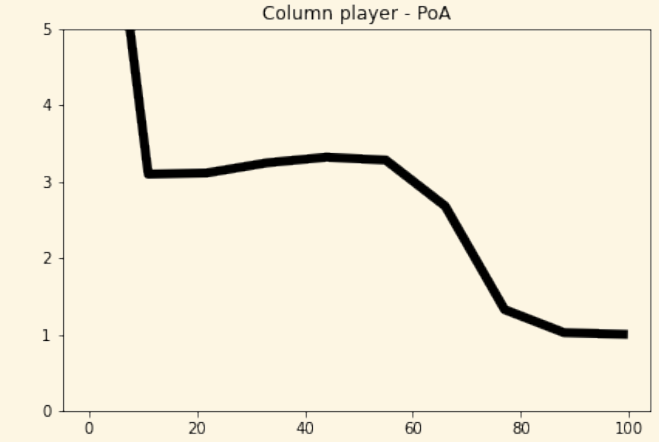
Inefficiencies can be learned and emerge naturally

Learning algorithms - Asymmetric replicator dynamics









Targeted incentivisation of behaviours can help escape learned inefficiencies

Thank you!

"Inefficiencies can be learned and emerge naturally"

"Targeted incentivisation of behaviours can help escape learned inefficiencies"

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@11michalis11
https://github.com/11michalis11/AmbulanceDecisionGame