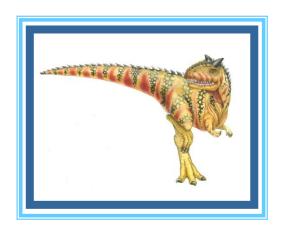
Chapter 10: File System

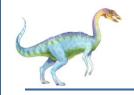




Chapter 10: File System

- n File Concept
- n Access Methods
- n Disk and Directory Structure
- n File-System Mounting
- n File Sharing
- n Protection

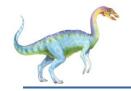




Objectives

- n To explain the function of file systems
- n To describe the interfaces to file systems
- n To discuss file-system design tradeoffs, including access methods, file sharing, file locking, and directory structures
- n To explore file-system protection





File Concept

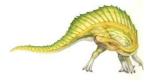
- n Contiguous logical address space
- n Types:
 - Data
 - numeric
 - character
 - binary
 - Program
- n Contents defined by file's creator
 - Many types
 - Consider text file, source file, executable file

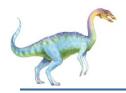




File Attributes

- Name only information kept in human-readable form
- n Identifier unique tag (number) identifies file within file system
- n Type needed for systems that support different types
- n Location pointer to file location on device
- n Size current file size
- n Protection controls who can do reading, writing, executing
- n Time, date, and user identification data for protection, security, and usage monitoring
- n Information about files are kept in the directory structure, which is maintained on the disk
- n Many variations, including extended file attributes such as file checksum
- n Information kept in the directory structure

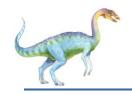




File info Window on Mac OS X

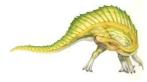


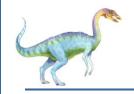




File Operations

- n File is an abstract data type
- n Create
- n Write at write pointer location
- n Read at read pointer location
- n Reposition within file seek
- n Delete
- n Truncate
- n Open (F_i) search the directory structure on disk for entry F_i , and move the content of entry to memory
- n Close (F_i) move the content of entry F_i in memory to directory structure on disk





Open Files

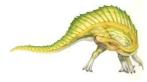
- n Several pieces of data are needed to manage open files:
 - Open-file table: tracks open files
 - File pointer: pointer to last read/write location, per process that has the file open
 - File-open count: counter of number of times a file is open to allow removal of data from open-file table when last processes closes it
 - Disk location of the file: cache of data access information
 - Access rights: per-process access mode information

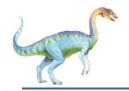




Open File Locking

- n Provided by some operating systems and file systems
 - Similar to reader-writer locks
 - Shared lock similar to reader lock several processes can acquire concurrently
 - Exclusive lock similar to writer lock
- n Mediates access to a file
- n Mandatory or advisory:
 - Mandatory access is denied depending on locks held and requested
 - Advisory processes can find status of locks and decide what to do





File Locking Example – Java API

```
import java.io.*;
import java.nio.channels.*;
public class LockingExample {
    public static final boolean EXCLUSIVE = false;
    public static final boolean SHARED = true;
    public static void main(String arsg[]) throws IOException {
            FileLock sharedLock = null;
           FileLock exclusiveLock = null;
           try {
                        RandomAccessFile raf = new RandomAccessFile("file.txt", "rw");
                       // get the channel for the file
                        FileChannel ch = raf.getChannel();
                       // this locks the first half of the file - exclusive
                       exclusiveLock = ch.lock(0, raf.length()/2, EXCLUSIVE);
                       /** Now modify the data . . . */
                       // release the lock
                       exclusiveLock.release();
```





File Locking Example – Java API (Cont.)

```
// this locks the second half of the file - shared
          sharedLock = ch.lock(raf.length()/2+1, raf.length(),
                                SHARED);
          /** Now read the data . . . */
          // release the lock
          sharedLock.release();
} catch (java.io.IOException ioe) {
          System.err.println(ioe);
}finally {
          if (exclusiveLock != null)
          exclusiveLock.release();
          if (sharedLock != null)
          sharedLock.release();
```

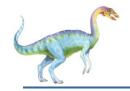




File Types – Name, Extension

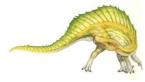
file type	usual extension	function	
executable	exe, com, bin or none	ready-to-run machine- language program	
object	obj, o	compiled, machine language, not linked	
source code	c, cc, java, pas, asm, a	source code in various languages	
batch	bat, sh	commands to the command interpreter	
text	txt, doc	textual data, documents	
word processor	wp, tex, rtf, doc	various word-processor formats	
library	lib, a, so, dll	libraries of routines for programmers	
print or view	ps, pdf, jpg	ASCII or binary file in a format for printing or viewing	
archive	arc, zip, tar	related files grouped into one file, sometimes com- pressed, for archiving or storage	
multimedia	mpeg, mov, rm, mp3, avi	binary file containing audio or A/V information	





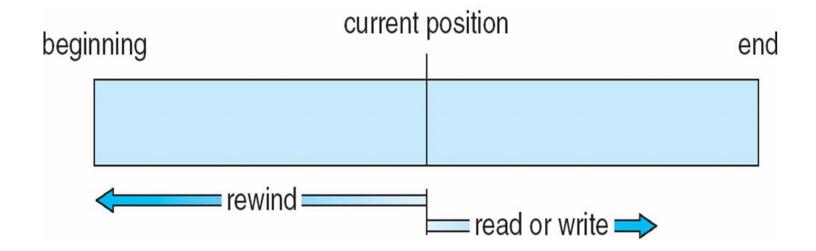
File Structure

- None sequence of words, bytes
- n Simple record structure
 - Lines
 - Fixed length
 - Variable length
- n Complex Structures
 - Formatted document
 - Relocatable load file
- n Can simulate last two with first method by inserting appropriate control characters
- n Who decides:
 - Operating system
 - Program

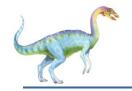




Sequential-access File







Access Methods

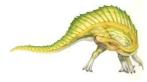
n Sequential Access

```
read next
write next
reset
no read after last write
(rewrite)
```

n Direct Access – file is fixed length logical records

n = relative block number

- n Relative block numbers allow OS to decide where file should be placed
 - See allocation problem in Ch 11





Simulation of Sequential Access on Direct-access File

sequential access	implementation for direct access	
reset	cp = 0;	
read next	read cp; $cp = cp + 1$;	
write next	write cp ; cp = cp + 1;	





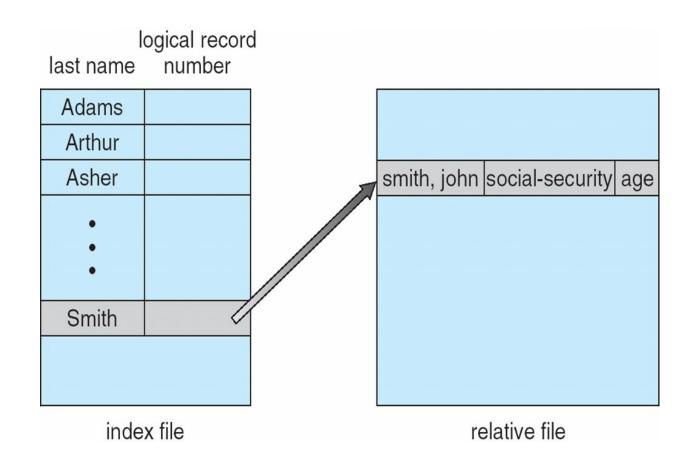
Other Access Methods

- n Can be built on top of base methods
- n General involve creation of an index for the file
- n Keep index in memory for fast determination of location of data to be operated on (consider UPC code plus record of data about that item)
- n If too large, index (in memory) of the index (on disk)
- n IBM indexed sequential-access method (ISAM)
 - Small master index, points to disk blocks of secondary index
 - File kept sorted on a defined key
 - All done by the OS
- vMS operating system provides index and relative files as another example (see next slide)

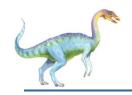




Example of Index and Relative Files

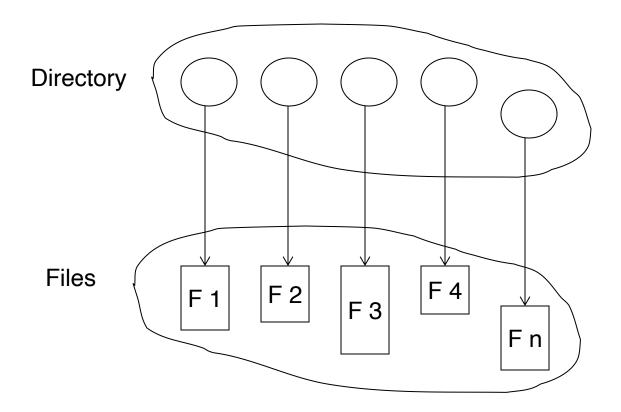




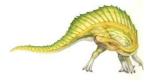


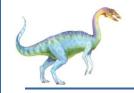
Directory Structure

n A collection of nodes containing information about all files



Both the directory structure and the files reside on disk





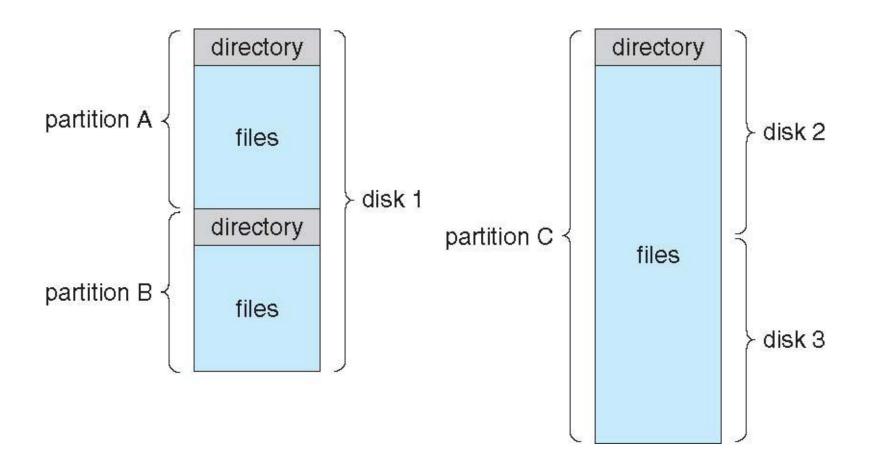
Disk Structure

- n Disk can be subdivided into partitions
- n Disks or partitions can be RAID protected against failure
- n Disk or partition can be used raw without a file system, or formatted with a file system
- n Partitions also known as minidisks, slices
- n Entity containing file system known as a volume
- n Each volume containing file system also tracks that file system's indevice directory or volume table of contents
- n As well as general-purpose file systems there are many specialpurpose file systems, frequently all within the same operating system or computer

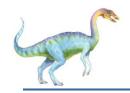




A Typical File-system Organization

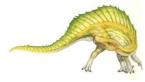






Types of File Systems

- n We mostly talk of general-purpose file systems
- n But systems frequently have many file systems, some general- and some special- purpose
- n Consider Solaris has
 - tmpfs memory-based volatile FS for fast, temporary I/O
 - objfs interface into kernel memory to get kernel symbols for debugging
 - ctfs contract file system for managing daemons
 - I lofs loopback file system allows one FS to be accessed in place of another
 - procfs kernel interface to process structures
 - ufs, zfs general purpose file systems





Operations Performed on Directory

- n Search for a file
- n Create a file
- n Delete a file
- n List a directory
- n Rename a file
- n Traverse the file system

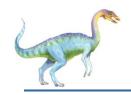




Organize the Directory (Logically) to Obtain

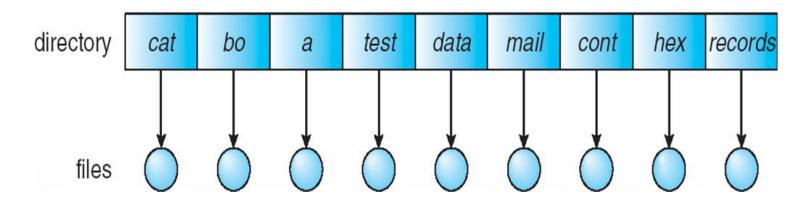
- n Efficiency locating a file quickly
- Naming convenient to users
 - I Two users can have same name for different files
 - The same file can have several different names
- n Grouping logical grouping of files by properties, (e.g., all Java programs, all games, ...)





Single-Level Directory

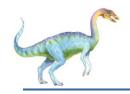
n A single directory for all users



Naming problem

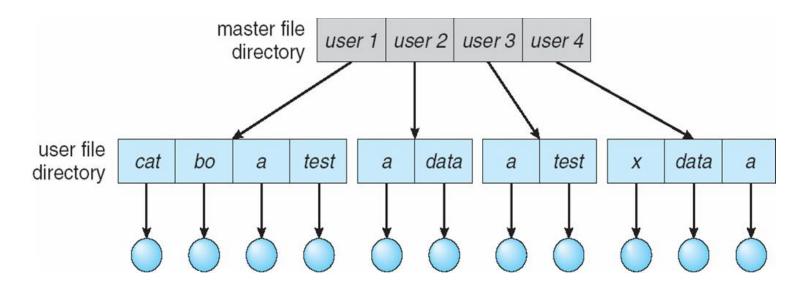
Grouping problem



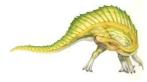


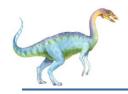
Two-Level Directory

n Separate directory for each user

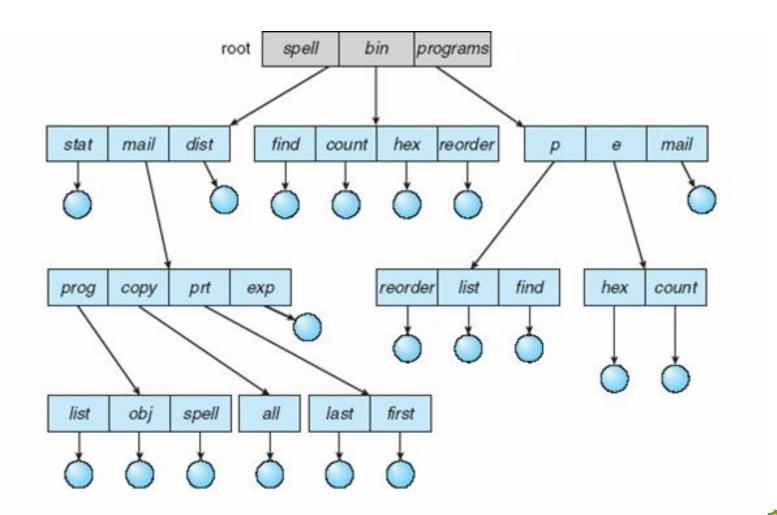


- n Path name
- n Can have the same file name for different user
- n Efficient searching
- No grouping capability

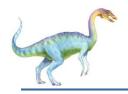




Tree-Structured Directories



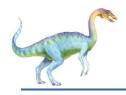




Tree-Structured Directories (Cont.)

- n Efficient searching
- n Grouping Capability
- n Current directory (working directory)
 - cd /spell/mail/prog
 - type list





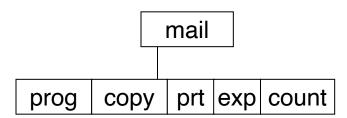
Tree-Structured Directories (Cont)

- n Absolute or relative path name
- n Creating a new file is done in current directory
- n Delete a file

n Creating a new subdirectory is done in current directory

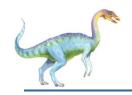
Example: if in current directory /mail

mkdir count



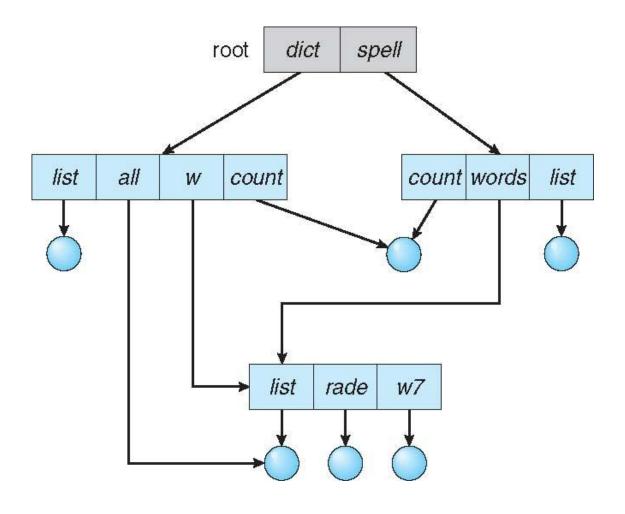
Deleting "mail" ⇒ deleting the entire subtree rooted by "mail"





Acyclic-Graph Directories

n Have shared subdirectories and files

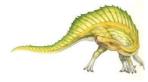






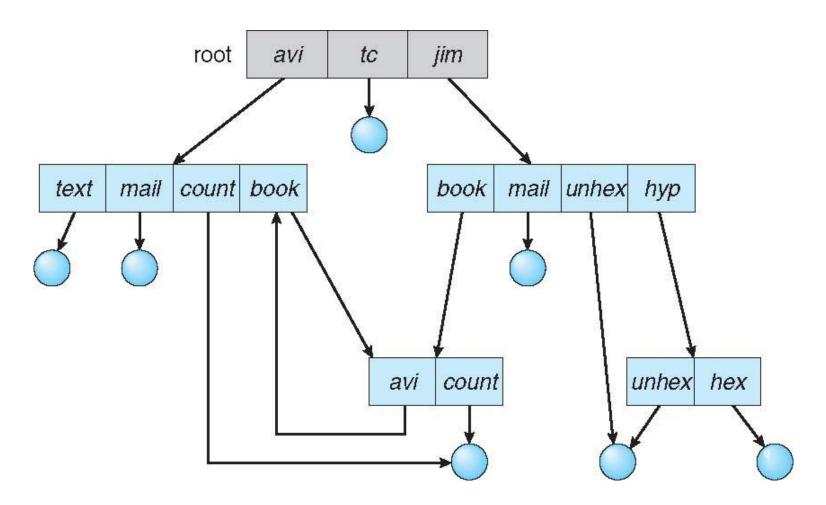
Acyclic-Graph Directories (Cont.)

- n Two different names (aliasing)
- n If dict deletes count ⇒ dangling pointer Solutions:
 - Backpointers, so we can delete all pointers Variable size records a problem
 - Backpointers using a daisy chain organization
 - Entry-hold-count solution
- n New directory entry type
 - Link another name (pointer) to an existing file
 - Resolve the link follow pointer to locate the file

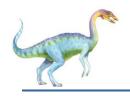




General Graph Directory







General Graph Directory (Cont.)

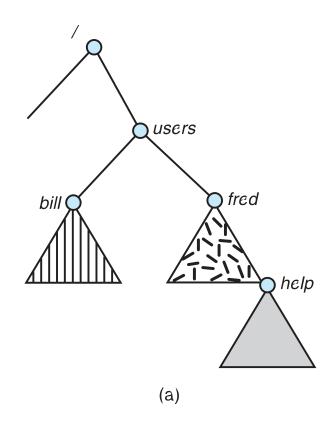
- n How do we guarantee no cycles?
 - Allow only links to file not subdirectories
 - I Garbage collection
 - Every time a new link is added use a cycle detection algorithm to determine whether it is OK

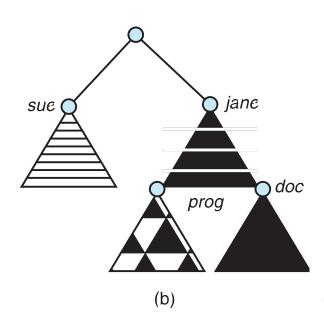


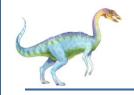


File System Mounting

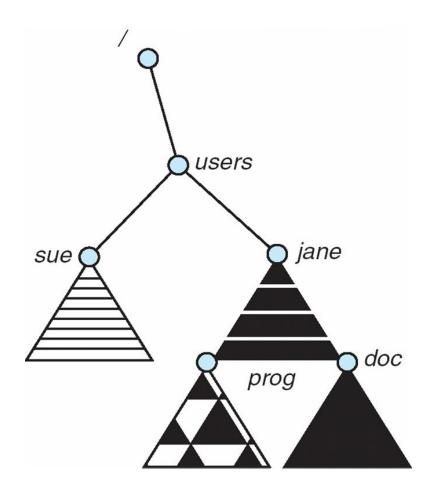
- n A file system must be mounted before it can be accessed
- n A unmounted file system is mounted at a mount point



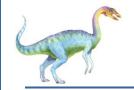




Mount Point



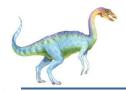




File Sharing

- n Sharing of files on multi-user systems is desirable
- n Sharing may be done through a protection scheme
- n On distributed systems, files may be shared across a network
- Network File System (NFS) is a common distributed file-sharing method
- n If multi-user system
 - User IDs identify users, allowing permissions and protections to be peruser
 - Group IDs allow users to be in groups, permitting group access rights
 - Owner of a file / directory
 - Group of a file / directory





File Sharing – Remote File Systems

- n Uses networking to allow file system access between systems
 - Manually via programs like FTP
 - Automatically, seamlessly using distributed file systems
 - Semi automatically via the world wide web
- n Client-server model allows clients to mount remote file systems from servers
 - Server can serve multiple clients
 - Client and user-on-client identification is insecure or complicated
 - NFS is standard UNIX client-server file sharing protocol
 - CIFS is standard Windows protocol
 - Standard operating system file calls are translated into remote calls
- n Distributed Information Systems (distributed naming services) such as LDAP, DNS, NIS, Active Directory implement unified access to information needed for remote computing



File Sharing – Failure Modes

- n All file systems have failure modes
 - For example corruption of directory structures or other non-user data, called metadata
- n Remote file systems add new failure modes, due to network failure, server failure
- n Recovery from failure can involve state information about status of each remote request
- n Stateless protocols such as NFS v3 include all information in each request, allowing easy recovery but less security



- n Specify how multiple users are to access a shared file simultaneously
 - Similar to Ch 6 process synchronization algorithms
 - Tend to be less complex due to disk I/O and network latency (for remote file systems
 - Andrew File System (AFS) implemented complex remote file sharing semantics
 - Unix file system (UFS) implements:
 - Writes to an open file visible immediately to other users of the same open file
 - Sharing file pointer to allow multiple users to read and write concurrently
 - AFS has session semantics
 - Writes only visible to sessions starting after the file is closed





Protection

- n File owner/creator should be able to control:
 - what can be done
 - by whom
- n Types of access
 - Read
 - Write
 - Execute
 - Append
 - Delete
 - l List



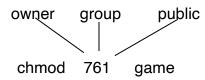


Access Lists and Groups

- Mode of access: read, write, execute
- n Three classes of users on Unix / Linux

			RWX
a) owner access	7	\Rightarrow	111
,			RWX
b) group access	6	\Rightarrow	110
, 5			RWX
c) public access	1	\Rightarrow	0 0 1

- n Ask manager to create a group (unique name), say G, and add some users to the group.
- n For a particular file (say game) or subdirectory, define an appropriate access.

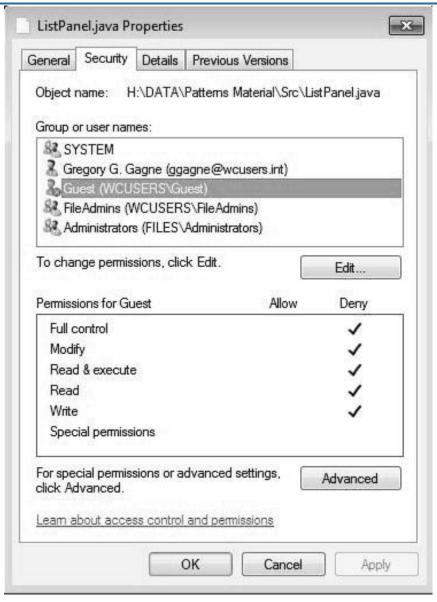


Attach a group to a file

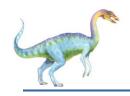
chgrp G game











A Sample UNIX Directory Listing

-rw-rw-r	1 pbg	staff	31200	Sep 3 08:30	intro.ps
drwx	5 pbg	staff	512	Jul 8 09.33	private/
drwxrwxr-x	2 pbg	staff	512	Jul 8 09:35	doc/
drwxrwx	2 pbg	student	512	Aug 3 14:13	student-proj/
-rw-rr	1 pbg	staff	9423	Feb 24 2003	program.c
-rwxr-xr-x	1 pbg	staff	20471	Feb 24 2003	program
drwxxx	4 pbg	faculty	512	Jul 31 10:31	lib/
drwx	3 pbg	staff	1024	Aug 29 06:52	mail/
drwxrwxrwx	3 pbg	staff	512	Jul 8 09:35	test/

