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一. 数论

1.阶乘最后非零位

//求阶乘最后非零位,复杂度 O(nlogn) //返回该位,n 以字符串方式传入 #include <string.h> #define MAXN 10000

```
int lastdigit(char* buf){
    const int mod[20]=\{1,1,2,6,4,2,2,4,2,8,4,4,8,4,6,8,8,6,8,2\};
    int len=strlen(buf),a[MAXN],i,c,ret=1;
    if (len==1)
       return mod[buf[0]-'0'];
    for (i=0;i<len;i++)</pre>
       a[i]=buf[len-1-i]-'0';
    for (;len;len-=!a[len-1]){
       ret=ret*mod[a[1]%2*10+a[0]]%5;
       for (c=0,i=len-1;i>=0;i--)
           c=c*10+a[i],a[i]=c/5,c%=5;
    }
    return ret+ret%2*5;
}
                                  2. 模线性方程(组)
#ifdef WIN32
typedef __int64 i64;
typedef long long i64;
#endif
//扩展 Euclid 求解 gcd(a,b)=ax+by
int ext_gcd(int a,int b,int& x,int& y){
    int t,ret;
    if (!b){
       x=1, y=0;
       return a;
    }
    ret=ext_gcd(b,a%b,x,y);
    t=x, x=y, y=t-a/b*y;
    return ret;
}
//计算 m^a, O(loga), 本身没什么用, 注意这个按位处理的方法:-P
int exponent(int m,int a){
    int ret=1;
    for (;a;a>>=1,m*=m)
       if (a&1)
           ret*=m;
    return ret;
}
//计算幂取模 a^b mod n, O(logb)
int modular_exponent(int a,int b,int n){ //a^b mod n
    int ret=1;
```

```
for (;b;b>>=1,a=(int)((i64)a)*a%n)
       if (b&1)
           ret=(int)((i64)ret)*a%n;
    return ret;
}
//求解模线性方程 ax=b (mod n)
//返回解的个数,解保存在 sol[]中
//要求 n>0,解的范围 0..n-1
int modular_linear(int a,int b,int n,int* sol){
    int d,e,x,y,i;
   d=ext_gcd(a,n,x,y);
   if (b%d)
       return 0;
   e=(x*(b/d)%n+n)%n;
    for (i=0; i< d; i++)
       sol[i]=(e+i*(n/d))%n;
    return d;
}
//求解模线性方程组(中国余数定理)
// x = b[0] \pmod{w[0]}
// x = b[1] \pmod{w[1]}
// ...
// x = b[k-1] \pmod{w[k-1]}
//要求 w[i]>0,w[i]与 w[j]互质,解的范围 1..n,n=w[0]*w[1]*...*w[k-1]
int modular_linear_system(int b[],int w[],int k){
    int d,x,y,a=0,m,n=1,i;
    for (i=0;i<k;i++)
       n*=w[i];
    for (i=0; i< k; i++){
       m=n/w[i];
       d=ext_gcd(w[i],m,x,y);
       a=(a+y*m*b[i])%n;
   }
    return (a+n)%n;
}
                                      3. 素数表
//用素数表判定素数,先调用 initprime
int plist[10000],pcount=0;
int prime(int n){
   int i;
    if ((n!=2\&\&!(n%2))||(n!=3\&\&!(n%3))||(n!=5\&\&!(n%5))||(n!=7\&\&!(n%7)))
        return 0;
```

```
for (i=0;plist[i]*plist[i]<=n;i++)</pre>
        if (!(n%plist[i]))
            return 0;
    return n>1;
}
void initprime(){
    int i;
    for (plist[pcount++]=2,i=3;i<50000;i++)
        if (prime(i))
            plist[pcount++]=i;
}
                            4. 素数随机判定(miller_rabin)
//miller rabin
//判断自然数 n 是否为素数
//time 越高失败概率越低,一般取 10 到 50
#include <stdlib.h>
#ifdef WIN32
typedef __int64 i64;
#else
typedef long long i64;
#endif
int modular_exponent(int a,int b,int n){ //a^b mod n
    int ret;
    for (;b;b>>=1,a=(int)((i64)a)*a%n)
        if (b&1)
            ret=(int)((i64)ret)*a%n;
    return ret;
}
// Carmicheal number: 561,41041,825265,321197185
int miller_rabin(int n,int time=10){
    if (n=1||(n!=2\&\&!(n%2))||(n!=3\&\&!(n%3))||(n!=5\&\&!(n%5))||(n!=7\&\&!(n%7)))
        return 0;
    while (time--)
        i f
(\mathsf{modular\_exponent}(((\mathsf{rand}()\&0x7fff<<16)+\mathsf{rand}()\&0x7fff+\mathsf{rand}()\&0x7fff)\%(\mathsf{n}-1)+1,\mathsf{n}-1,\mathsf{n})\,!=1)
            return 0;
    return 1;
}
                                       5. 质因数分解
//分解质因数
//prime_factor()传入 n,返回不同质因数的个数
```

```
//f 存放质因数, nf 存放对应质因数的个数
//先调用 initprime(), 其中第二个 initprime()更快
#include<iostream>
#include<cstdio>
#include<cmath>
using namespace std;
#define MAXN 2001000
#define PSIZE 100000
int plist[PSIZE], pcount=0;
int prime(int n){
    int i;
    if ((n!=2\&\&!(n%2))||(n!=3\&\&!(n%3))||(n!=5\&\&!(n%5))||(n!=7\&\&!(n%7)))
        return 0;
    for (i=0;plist[i]*plist[i]<=n;++i)</pre>
        if (!(n%plist[i]))
           return 0;
    return n>1;
}
void initprime(){
    int i;
   for (plist[pcount++]=2,i=3;i<100000;++i)</pre>
       if (prime(i))
           plist[pcount++]=i;
}
int prime_factor(int n, int* f, int *nf) {
   int cnt = 0;
    int n2 = sqrt((double)n);
    for(int i = 0; n > 1 \&\& plist[i] <= n2; ++i)
        if (n % plist[i] == 0) {
           for (nf[cnt] = 0; n \% plist[i] == 0; ++nf[cnt], n /= plist[i]);
           f[cnt++] = plist[i];
    if (n > 1) nf[cnt] = 1, f[cnt++] = n;
    return cnt;
}
/*
//产生 MAXN 以内的所有素数
//note:2863311530 就是 10101010101010101010101010101010
//给所有 2 的倍数赋初值
#include <cmath>
#include <iostream>
using namespace std;
#define MAXN 10000000
```

```
unsigned int plist[6000000],pcount;
unsigned int isprime[(MAXN>>5)+1];
#define setbitzero(a) (isprime[(a)>>5]&=(\sim(1<<((a)&31))))
#define setbitone(a) (isprime[(a)>>5]|=(1<<((a)&31)))
#define ISPRIME(a) (isprime[(a) >> 5]&(1 << ((a)&31)))
void initprime(){
   int i,j,m;
   int t=(MAXN>>5)+1;
   for(i=0;i<t;++i)isprime[i]=2863311530;</pre>
   plist[0]=2; setbitone(2); setbitzero(1);
   m=(int)sqrt(MAXN);
   for(pcount=1,i=3;i<=m;i+=2)</pre>
       if(ISPRIME(i))
          for(plist[pcount++]=i,j=i<<1;j<=MAXN;j+=i)</pre>
              setbitzero(j);
   if(!(i&1))++i;
   for(;i<=MAXN;i+=2)if(ISPRIME(i))plist[pcount++]=i;</pre>
}
                                 6. 最大公约数欧拉函数
int gcd(int a,int b){
    return b?gcd(b,a%b):a;
}
inline int lcm(int a,int b){
    return a/gcd(a,b)*b;
}
//求 1..n-1 中与 n 互质的数的个数
int eular(int n){
    int ret=1,i;
    for (i=2;i*i<=n;i++)
        if (n\%i==0){
            n/=i,ret*=i-1;
            while (n\%i==0)
                n/=i,ret*=i;
        }
    if (n>1)
        ret*=n-1;
    return ret;
}
```

二.图论_匹配

1. 二分图最大匹配(hungary 邻接表形式)

```
//二分图最大匹配, hungary 算法, 邻接表形式, 复杂度 O(m*e)
//返回最大匹配数,传入二分图大小 m,n 和邻接表 list(只需一边)
//match1, match2 返回一个最大匹配, 未匹配顶点 match 值为-1
#include <string.h>
#define MAXN 310
#define _clr(x) memset(x,0xff,sizeof(int)*MAXN)
struct edge_t{
   int from, to;
   edge_t* next;
};
int hungary(int m,int n,edge_t* list[],int* match1,int* match2){
   int s[MAXN],t[MAXN],p,q,ret=0,i,j,k;edge t* e;
   for (_clr(match1),_clr(match2),i=0;i<m;ret+=(match1[i++]>=0))
       for (_clr(t), s[p=q=0]=i; p <= q\&match1[i] < 0; p++)
           for (e=list[k=s[p]];e&&match1[i]<0;e=e->next)
               if (t[j=e->to]<0){
                  s[++q]=match2[j],t[j]=k;
                  if (s[q]<0)
                      for (p=j;p>=0;j=p)
                          match2[j]=k=t[j],p=match1[k],match1[k]=j;
               }
   return ret;
}
               2. 二分图最大匹配(hungary 邻接表形式,邻接阵接口)
//二分图最大匹配, hungary 算法, 邻接表形式, 邻接阵接口, 复杂度 0(m*e)s
//返回最大匹配数,传入二分图大小 m,n 和邻接阵
//match1, match2 返回一个最大匹配, 未匹配顶点 match 值为-1
#include <string.h>
#include <vector>
#define MAXN 310
#define _clr(x) memset(x,0xff,sizeof(int)*MAXN)
int hungary(int m,int n,int mat[][MAXN],int* match1,int* match2){
   int s[MAXN],t[MAXN],p,q,ret=0,i,j,k,r;
   vector<int> e[MAXN];
   //生成邻接表(只需一边)
   for(i=0;i<m;++i)</pre>
       for(j=0;j<n;++j)
```

```
if (mat[i][j]) e[i].push_back(j);
   for (_clr(match1),_clr(match2),i=0;i<m;ret+=(match1[i++]>=0))
       for (clr(t), s[p=q=0]=i; p <= q\&match1[i] < 0; p++)
           for(r=0,k=s[p];r<e[k].size()&&match1[i]<0;++r)
               if (t[j=e[k][r]]<0){
                  s[++q]=match2[j],t[j]=k;
                  if (s[q]<0)
                      for (p=j;p>=0;j=p)
                          match2[j]=k=t[j],p=match1[k],match1[k]=j;
               }
   return ret;
}
                     3. 二分图最大匹配(hungary 邻接阵形式)
//二分图最大匹配, hungary 算法, 邻接阵形式, 复杂度 0(m*m*n)
//返回最大匹配数,传入二分图大小 m,n 和邻接阵 mat,非零元素表示有边
//match1, match2返回一个最大匹配, 未匹配顶点 match 值为-1
#include <string.h>
#define MAXN 310
#define _clr(x) memset(x,0xff,sizeof(int)*MAXN)
int hungary(int m,int n,int mat[][MAXN],int* match1,int* match2){
   int s[MAXN],t[MAXN],p,q,ret=0,i,j,k;
   for (_clr(match1),_clr(match2),i=0;i<m;ret+=(match1[i++]>=0))
       for (_clr(t), s[p=q=0]=i; p <= q\&match1[i] < 0; p++)
           for (k=s[p],j=0;j<n&&match1[i]<0;j++)</pre>
               if (mat[k][j]\&\&t[j]<0){
                  s[++q]=match2[i],t[i]=k;
                  if (s[q]<0)
                      for (p=j;p>=0;j=p)
                          match2[j]=k=t[j],p=match1[k],match1[k]=j;
               }
   return ret;
}
                     4. 二分图最大匹配(hungary 正向表形式)
//二分图最大匹配, hungary 算法, 正向表形式, 复杂度 0 (m*e)
//返回最大匹配数,传入二分图大小 m,n 和正向表 list,buf(只需一边)
//match1, match2返回一个最大匹配, 未匹配顶点 match 值为-1
#include <string.h>
#define MAXN 310
#define _clr(x) memset(x,0xff,sizeof(int)*MAXN)
int hungary(int m,int n,int* list,int* buf,int* match1,int* match2){
   int s[MAXN],t[MAXN],p,q,ret=0,i,j,k,l;
   for (_clr(match1),_clr(match2),i=0;i<m;ret+=(match1[i++]>=0))
```

```
for (_clr(t), s[p=q=0]=i; p <= q\&match1[i] < 0; p++)
            for (l=list[k=s[p]]; l<list[k+1]&&match1[i]<0; l++)</pre>
                if (t[j=buf[l]]<0){</pre>
                   s[++q]=match2[j],t[j]=k;
                   if (s[q]<0)
                       for (p=j;p>=0;j=p)
                           match2[j]=k=t[j],p=match1[k],match1[k]=j;
               }
    return ret;
}
                   5. 二分图最佳匹配(kuhn_munkras 邻接阵形式)
//二分图最佳匹配,kuhn munkras 算法,邻接阵形式,复杂度 0(m*m*n)
//返回最佳匹配值,传入二分图大小 m,n 和邻接阵 mat,表示权值
//match1, match2返回一个最佳匹配, 未匹配顶点 match 值为-1
//一定注意 m<=n, 否则循环无法终止
//最小权匹配可将权值取相反数
#include <string.h>
#define MAXN 310
#define inf 1000000000
#define clr(x) memset(x,0xff,sizeof(int)*n)
int kuhn_munkras(int m,int n,int mat[][MAXN],int* match1,int* match2){
    int s[MAXN],t[MAXN],l1[MAXN],l2[MAXN],p,q,ret=0,i,j,k;
    for (i=0; i < m; i++)
        for (l1[i]=-inf, j=0; j<n; j++)
            l1[i]=mat[i][j]>l1[i]?mat[i][j]:l1[i];
    for (i=0; i< n; l2[i++]=0);
    for (_clr(match1),_clr(match2),i=0;i<m;i++){</pre>
        for (_clr(t), s[p=q=0]=i; p <= q\&match1[i] < 0; p++)
            for (k=s[p],j=0;j<n&&match1[i]<0;j++)</pre>
                if (l1[k]+l2[j]==mat[k][j]&&t[j]<0){
                   s[++q]=match2[j],t[j]=k;
                   if (s[q]<0)
                        for (p=j;p>=0;j=p)
                           match2[j]=k=t[j],p=match1[k],match1[k]=j;
                }
        if (match1[i]<0){</pre>
            for (i--,p=inf,k=0;k<=q;k++)
                for (j=0; j< n; j++)
                   if (t[j]<0\&l1[s[k]]+l2[j]-mat[s[k]][j]<p)
                       p=l1[s[k]]+l2[j]-mat[s[k]][j];
            for (j=0; j< n; l2[j]+=t[j]<0?0:p, j++);
            for (k=0; k \le q; l1[s[k++]] -= p);
       }
    }
```

```
for (i=0;i<m;i++)
        ret+=mat[i][match1[i]];
    return ret;
}
                             6. 一般图匹配(邻接表形式)
//一般图最大匹配,邻接表形式,复杂度 0(n*e)
//返回匹配顶点对数, match 返回匹配, 未匹配顶点 match 值为-1
//传入图的顶点数 n 和邻接表 list
#define MAXN 100
struct edge_t{
    int from, to;
   edge_t* next;
};
int aug(int n,edge_t* list[],int* match,int* v,int now){
    int t,ret=0;edge_t* e;
   v[now]=1;
    for (e=list[now];e;e=e->next)
       if (!v[t=e->to]){
           if (match[t]<0)</pre>
               match[now]=t,match[t]=now,ret=1;
           else{
               v[t]=1;
               if (aug(n,list,match,v,match[t]))
                   match[now]=t,match[t]=now,ret=1;
               v[t]=0;
           }
           if (ret)
               break;
       }
    v[now]=0;
    return ret;
}
int graph_match(int n,edge_t* list[],int* match){
    int v[MAXN],i,j;
    for (i=0;i<n;i++)
   v[i]=0,match[i]=-1;
    for (i=0, j=n; i< n\&\&j>=2;)
       if (match[i]<0&&aug(n,list,match,v,i))</pre>
           i=0, j-=2;
       else
           i++;
    for (i=j=0;i< n;i++)
       j+=(match[i]>=0);
```

```
return j/2;
}
```

7. 一般图匹配(邻接表形式,邻接阵接口)

```
//一般图最大匹配,邻接表形式,复杂度 0(n*e)
//返回匹配顶点对数, match 返回匹配, 未匹配顶点 match 值为-1
//传入图的顶点数 n 和邻接表 list
#include <vector>
#define MAXN 100
int aug(int n,vector<int> list[],int* match,int* v,int now){
    int t,ret=0,r;
   v[now]=1;
// for (e=list[now];e;e=e->next)
    for (r=0;r<list[now].size();++r)</pre>
        if (!v[t=list[now][r]]){
            if (match[t]<0)</pre>
               match[now]=t,match[t]=now,ret=1;
            else{
               v[t]=1;
                if (aug(n,list,match,v,match[t]))
                    match[now]=t,match[t]=now,ret=1;
               v[t]=0;
            }
            if (ret)
               break;
       }
    v[now]=0;
    return ret;
}
int graph_match(int n,int mat[][MAXN],int* match){
    int v[MAXN],i,j;
   vector<int> list[MAXN];
    for (i=0;i<n;i++)
        for (j=0; j< n; j++)
            if (mat[i][j]) list[i].push_back(j);
    for (i=0;i<n;i++)
    v[i]=0, match[i]=-1;
    for (i=0, j=n; i< n\&\&j>=2;)
        if (match[i]<0&&aug(n,list,match,v,i))</pre>
            i=0, j-=2;
        else
            i++;
    for (i=j=0;i< n;i++)
        j+=(match[i]>=0);
```

```
return j/2;
}
                             8. 一般图匹配(邻接阵形式)
//一般图最大匹配,邻接阵形式,复杂度 0(n^3)
//返回匹配顶点对数, match 返回匹配, 未匹配顶点 match 值为-1
//传入图的顶点数 n 和邻接阵 mat
#define MAXN 100
int aug(int n,int mat[][MAXN],int* match,int* v,int now){
   int i,ret=0;
   v[now]=1;
   for (i=0;i<n;i++)
       if (!v[i]&&mat[now][i]){
           if (match[i]<0)</pre>
               match[now]=i,match[i]=now,ret=1;
           else{
              v[i]=1;
               if (aug(n,mat,match,v,match[i]))
                  match[now]=i,match[i]=now,ret=1;
              v[i]=0;
           }
           if (ret)
              break;
       }
   v[now]=0;
   return ret;
}
int graph_match(int n,int mat[][MAXN],int* match){
   int v[MAXN],i,j;
   for (i=0;i<n;i++)
   v[i]=0,match[i]=-1;
   for (i=0, j=n; i< n\&\&j>=2;)
       if (match[i]<0&&aug(n,mat,match,v,i))</pre>
           i=0, j-=2;
       else
           i++;
   for (i=j=0;i< n;i++)
       j+=(match[i]>=0);
   return j/2;
}
                             9. 一般图匹配(正向表形式)
//一般图最大匹配,正向表形式,复杂度 0(n*e)
```

```
//传入图的顶点数 n 和正向表 list, buf
#define MAXN 100
int aug(int n,int* list,int* buf,int* match,int* v,int now){
    int i,t,ret=0;
    v[now]=1;
    for (i=list[now];i<list[now+1];i++)</pre>
        if (!v[t=buf[i]]){
            if (match[t]<0)</pre>
                match[now]=t,match[t]=now,ret=1;
            else{
                v[t]=1;
                if (aug(n,list,buf,match,v,match[t]))
                    match[now]=t,match[t]=now,ret=1;
                v[t]=0;
            }
            if (ret)
                break;
        }
    v[now]=0;
    return ret;
}
int graph_match(int n,int* list,int* buf,int* match){
    int v[MAXN],i,j;
    for (i=0;i< n;i++)
    v[i]=0,match[i]=-1;
    for (i=0, j=n; i< n\&\&j>=2;)
        if (match[i]<0&&aug(n,list,buf,match,v,i))</pre>
            i=0, j-=2;
        else
            i++;
    for (i=j=0;i< n;i++)
        j+=(match[i]>=0);
    return j/2;
}
```

三.图论_生成树

1. 最小生成树(kruskal 邻接表形式)

```
//无向图最小生成树, kruskal 算法, 邻接表形式, 复杂度 0 (mlogm) //返回最小生成树的长度, 传入图的大小 n 和邻接表 list //可更改边权的类型, edge[][2]返回树的构造, 用边集表示 //如果图不连通,则对各连通分支构造最小生成树, 返回总长度
```

```
#include <string.h>
#define MAXN 200
#define inf 1000000000
typedef double elem_t;
struct edge_t{
    int from, to;
    elem_t len;
    edge_t* next;
};
#define _ufind_run(x) for(;p[t=x];x=p[x],p[t]=(p[x]?p[x]:x))
#define _run_both _ufind_run(i);_ufind_run(j)
struct ufind{
    int p[MAXN],t;
    void init(){memset(p,0,sizeof(p));}
    void set_friend(int i,int j){_run_both;p[i]=(i==j?0:j);}
    int is_friend(int i,int j){_run_both;return i==j&&i;}
};
#define _cp(a,b) ((a).len<(b).len)</pre>
struct heap_t{int a,b;elem_t len;};
struct minheap{
    heap_t h[MAXN*MAXN];
    int n,p,c;
    void init(){n=0;}
    void ins(heap_t e){
        for (p=++n;p>1&\&_cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
        h[p]=e;
    int del(heap_t& e){
        if (!n) return 0;
(e=h[p=1],c=2;c<n\&\&\_cp(h[c+=(c< n-1\&\&\_cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
        h[p]=h[n--];return 1;
    }
};
elem_t kruskal(int n,edge_t* list[],int edge[][2]){
    ufind u; minheap h;
    edge_t* t;heap_t e;
    elem_t ret=0;int i,m=0;
    u.init(),h.init();
    for (i=0;i<n;i++)</pre>
        for (t=list[i];t;t=t->next)
            if (i<t->to)
                e.a=i,e.b=t->to,e.len=t->len,h.ins(e);
```

```
while (m<n-1&&h.del(e))
    if (!u.is_friend(e.a+1,e.b+1))
        edge[m][0]=e.a,edge[m][1]=e.b,ret+=e.len,u.set_friend(e.a+1,e.b+1);
    return ret;
}</pre>
```

2. 最小生成树(kruskal 正向表形式)

```
//无向图最小生成树, kruskal 算法, 正向表形式, 复杂度 O(mlogm)
//返回最小生成树的长度,传入图的大小n和正向表 list,buf
//可更改边权的类型,edge[][2]返回树的构造,用边集表示
//如果图不连通,则对各连通分支构造最小生成树,返回总长度
#include <string.h>
#define MAXN 200
#define inf 1000000000
typedef double elem_t;
struct edge_t{
   int to;
   elem_t len;
};
#define _ufind_run(x) for(;p[t=x];x=p[x],p[t]=(p[x]?p[x]:x))
#define _run_both _ufind_run(i);_ufind_run(j)
struct ufind{
   int p[MAXN],t;
   void init(){memset(p,0,sizeof(p));}
   void set friend(int i,int j){ run both;p[i]=(i==j?0:j);}
   int is_friend(int i,int j){_run_both;return i==j&&i;}
};
#define _{cp(a,b)} ((a).len<(b).len)
struct heap_t{int a,b;elem_t len;};
struct minheap{
   heap_t h[MAXN*MAXN];
   int n,p,c;
   void init(){n=0;}
   void ins(heap_t e){
       for (p=++n;p>1&\&_cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
       h[p]=e;
   }
   int del(heap t& e){
       if (!n) return 0;
       for
(e=h[p=1],c=2;c<n\&\&\_cp(h[c+=(c< n-1\&\&\_cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
       h[p]=h[n--];return 1;
   }
};
```

```
elem_t kruskal(int n,int* list,edge_t* buf,int edge[][2]){
   ufind u; minheap h;
   heap_t e;elem_t ret=0;
   int i,j,m=0;
   u.init(),h.init();
   for (i=0;i<n;i++)
       for (j=list[i];j<list[i+1];j++)</pre>
           if (i<buf[j].to)</pre>
               e.a=i,e.b=buf[j].to,e.len=buf[j].len,h.ins(e);
   while (m< n-1\&\&h.del(e))
       if (!u.is_friend(e.a+1,e.b+1))
           \verb|edge[m][0]=e.a, \verb|edge[m][1]=e.b, \verb|ret+=e.len, \verb|u.set_friend(e.a+1, e.b+1)|||
   return ret;
}
                  3. 最小生成树(prim+binary_heap 邻接表形式)
//无向图最小生成树,prim 算法+二分堆,邻接表形式,复杂度 O(mlogm)
//返回最小生成树的长度,传入图的大小n和邻接表 list
//可更改边权的类型,pre[]返回树的构造,用父结点表示,根节点(第一个)pre 值为-1
//必须保证图的连通的!
#define MAXN 200
#define inf 1000000000
typedef double elem_t;
struct edge_t{
   int from, to;
   elem_t len;
   edge_t* next;
};
#define _{cp(a,b)} ((a).d<(b).d)
struct heap_t{elem_t d;int v;};
struct heap{
   heap_t h[MAXN*MAXN];
   int n,p,c;
   void init(){n=0;}
   void ins(heap_t e){
       for (p=++n;p>1&\&_cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
       h[p]=e;
   }
   int del(heap_t& e){
       if (!n) return 0;
(e=h[p=1],c=2;c<n\&\&\_cp(h[c+=(c< n-1\&\&\_cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
       h[p]=h[n--];return 1;
   }
```

```
};
elem_t prim(int n,edge_t* list[],int* pre){
   heap h;
   elem_t min[MAXN],ret=0;
   edge_t* t;heap_t e;
   int v[MAXN],i;
   for (i=0;i<n;i++)
       min[i]=inf,v[i]=0,pre[i]=-1;
   h.init();e.v=0,e.d=0,h.ins(e);
   while (h.del(e))
       if (!v[e.v])
           for (v[e.v]=1, ret+=e.d, t=list[e.v]; t; t=t->next)
               if (!v[t->to]\&t->len<min[t->to])
                   pre[t->to]=t->from,min[e.v=t->to]=e.d=t->len,h.ins(e);
    return ret;
}
                  4. 最小生成树(prim+binary heap 正向表形式)
//无向图最小生成树,prim 算法+二分堆,正向表形式,复杂度 0(mlogm)
//返回最小生成树的长度,传入图的大小n和正向表 list,buf
//可更改边权的类型,pre[]返回树的构造,用父结点表示,根节点(第一个)pre 值为-1
//必须保证图的连通的!
#define MAXN 200
#define inf 1000000000
typedef double elem t;
struct edge_t{
   int to;
   elem_t len;
};
#define _{cp(a,b)} ((a).d<(b).d)
struct heap_t{elem_t d;int v;};
struct heap{
   heap t h[MAXN*MAXN];
   int n,p,c;
   void init(){n=0;}
   void ins(heap_t e){
       for (p=++n;p>1&\&\_cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
       h[p]=e;
   }
   int del(heap t& e){
       if (!n) return 0;
       for
(e=h[p=1], c=2; c<n\&\&\_cp(h[c+=(c<n-1\&\&\_cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
       h[p]=h[n--];return 1;
```

```
}
};
elem_t prim(int n,int* list,edge_t* buf,int* pre){
   heap h;heap_t e;
   elem_t min[MAXN],ret=0;
   int v[MAXN],i,j;
   for (i=0; i< n; i++)
       min[i]=inf,v[i]=0,pre[i]=-1;
   h.init();e.v=0,e.d=0,h.ins(e);
   while (h.del(e))
       if (!v[i=e.v])
           for (v[i]=1,ret+=e.d,j=list[i];j<list[i+1];j++)</pre>
               if (!v[buf[j].to]&&buf[j].len<min[buf[j].to])</pre>
                   pre[buf[j].to]=i,min[e.v=buf[j].to]=e.d=buf[j].len,h.ins(e);
    return ret;
}
                  5. 最小生成树(prim+mapped heap 邻接表形式)
//无向图最小生成树,prim 算法+映射二分堆,邻接表形式,复杂度 0(mlogn)
//返回最小生成树的长度,传入图的大小n和邻接表 list
//可更改边权的类型,pre[]返回树的构造,用父结点表示,根节点(第一个)pre 值为-1
//必须保证图的连通的!
#define MAXN 200
#define inf 1000000000
typedef double elem t;
struct edge_t{
   int from, to;
   elem_t len;
   edge_t* next;
};
#define _{cp(a,b)}((a)<(b))
struct heap{
   elem t h[MAXN+1];
   int ind[MAXN+1],map[MAXN+1],n,p,c;
   void init(){n=0;}
   void ins(int i,elem_t e){
       for (p=++n;p>1&\&\_cp(e,h[p>>1]);h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
       h[map[ind[p]=i]=p]=e;
   }
   int del(int i,elem t& e){
       i=map[i];if (i<1||i>n) return 0;
       for (e=h[p=i];p>1;h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
       for
(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<
```

```
=1);
       h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
   }
   int delmin(int& i,elem_t& e){
       if (n<1) return 0;i=ind[1];</pre>
(e=h[p=1], c=2; c<n&_cp(h[c+=(c<n-1&&_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c]
],p=c,c<<=1);
       h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
   }
};
elem_t prim(int n,edge_t* list[],int* pre){
   heap h;
   elem_t min[MAXN],ret=0,e;
   edge_t* t;
   int v[MAXN],i;
   for (h.init(),i=0;i<n;i++)
       min[i]=(i?inf:0),v[i]=0,pre[i]=-1,h.ins(i,min[i]);
   while (h.delmin(i,e))
       for (v[i]=1,ret+=e,t=list[i];t;t=t->next)
           if (!v[t->to]\&\&t->len<min[t->to])
               pre[t->to]=t->from,h.del(t->to,e),h.ins(t->to,min[t->to]=t->len);
   return ret;
}
                  6. 最小生成树(prim+mapped heap 正向表形式)
//无向图最小生成树,prim 算法+映射二分堆,正向表形式,复杂度 0(mlogn)
//返回最小生成树的长度,传入图的大小n和正向表 list,buf
//可更改边权的类型,pre[]返回树的构造,用父结点表示,根节点(第一个)pre 值为-1
//必须保证图的连通的!
#define MAXN 200
#define inf 1000000000
typedef double elem_t;
struct edge t{
   int to;
   elem_t len;
};
#define _{cp(a,b)}((a)<(b))
struct heap{
   elem t h[MAXN+1];
   int ind[MAXN+1],map[MAXN+1],n,p,c;
   void init(){n=0;}
   void ins(int i,elem_t e){
       for (p=++n;p>1&\&\_cp(e,h[p>>1]);h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
```

```
h[map[ind[p]=i]=p]=e;
          }
          int del(int i,elem t& e){
                     i=map[i];if (i<1||i>n) return 0;
                     for (e=h[p=i];p>1;h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[n]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[
=1);
                     h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
          }
          int delmin(int& i,elem_t& e){
                     if (n<1) return 0;i=ind[1];</pre>
(e=h[p=1],c=2;c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c]
],p=c,c<<=1);
                     h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
          }
};
elem_t prim(int n,int* list,edge_t* buf,int* pre){
          heap h;
          elem_t min[MAXN],ret=0,e;
          int v[MAXN],i,j;
          for (h.init(),i=0;i<n;i++)</pre>
                    min[i]=(i?inf:0),v[i]=0,pre[i]=-1,h.ins(i,min[i]);
          while (h.delmin(i,e))
                     for (v[i]=1,ret+=e,j=list[i];j<list[i+1];j++)</pre>
                               if (!v[buf[j].to]&&buf[j].len<min[buf[j].to])</pre>
          pre[buf[j].to]=i,h.del(buf[j].to,e),h.ins(buf[j].to,min[buf[j].to]=buf[j].len);
          return ret;
}
                                                                        7. 最小生成树(prim 邻接阵形式)
//无向图最小生成树,prim 算法,邻接阵形式,复杂度 0(n^2)
//返回最小生成树的长度,传入图的大小 n 和邻接阵 mat,不相邻点边权 inf
//可更改边权的类型,pre[]返回树的构造,用父结点表示,根节点(第一个)pre 值为-1
//必须保证图的连通的!
#define MAXN 200
#define inf 1000000000
typedef double elem_t;
elem_t prim(int n,elem_t mat[][MAXN],int* pre){
          elem_t min[MAXN],ret=0;
          int v[MAXN],i,j,k;
          for (i=0; i< n; i++)
```

```
min[i]=inf,v[i]=0,pre[i]=-1;
    for (\min[j=0]=0; j< n; j++){
        for (k=-1, i=0; i< n; i++)
            if (!v[i]&&(k==-1||min[i]<min[k]))
               k=i;
        for (v[k]=1,ret+=min[k],i=0;i<n;i++)</pre>
            if (!v[i]&&mat[k][i]<min[i])</pre>
               min[i]=mat[pre[i]=k][i];
    }
    return ret;
}
                              8. 最小树形图(邻接阵形式)
//多源最小树形图, edmonds 算法, 邻接阵形式, 复杂度 0(n^3)
//返回最小生成树的长度,构造失败返回负值
//传入图的大小 n 和邻接阵 mat, 不相邻点边权 inf
//可更改边权的类型,pre[]返回树的构造,用父结点表示
//传入时 pre[]数组清零,用-1标出源点
#include <string.h>
#define MAXN 120
#define inf 1000000000
typedef int elem_t;
elem_t edmonds(int n,elem_t mat[][MAXN*2],int* pre){
    elem_t ret=0;
    int c[MAXN*2][MAXN*2], l[MAXN*2], p[MAXN*2], m=n,t,i,j,k;
    for (i=0;i<n;l[i]=i,i++);
    do{
       memset(c,0,sizeof(c)),memset(p,0xff,sizeof(p));
        for (t=m,i=0;i<m;c[i][i]=1,i++);
        for (i=0;i<t;i++)
            if (l[i]==i&&pre[i]!=-1){
               for (j=0; j< m; j++)
                    if (l[j]==j\&\&i!=j\&\&mat[j][i]<inf\&\&(p[i]==-1||mat[j][i]<mat[p[i]][i])) \\
                       p[i]=j;
               if ((pre[i]=p[i])==-1)
                    return -1;
               if (c[i][p[i]]){
                   for (j=0;j<=m;mat[j][m]=mat[m][j]=inf,j++);</pre>
                   for (k=i;l[k]!=m;l[k]=m,k=p[k])
                       for (j=0; j< m; j++)
                           if (l[j]==j){
                               if (mat[j][k]-mat[p[k]][k]<mat[j][m])</pre>
                                   mat[j][m]=mat[j][k]-mat[p[k]][k];
                               if (mat[k][j]<mat[m][j])</pre>
                                   mat[m][j]=mat[k][j];
```

```
c[m][m]=1,l[m]=m,m++;
                }
                for (j=0; j< m; j++)
                    if (c[i][j])
                        for (k=p[i];k!=-1&\&l[k]==k;c[k][j]=1,k=p[k]);
            }
    }
   while (t<m);
    for (;m-->n;pre[k]=pre[m])
        for (i=0; i < m; i++)
            if (l[i]==m){
                for (j=0; j< m; j++)
                    if (pre[j]==m&&mat[i][j]==mat[m][j])
                         pre[i]=i;
                if (mat[pre[m]][m]==mat[pre[m]][i]-mat[pre[i]][i])
            }
    for (i=0;i<n;i++)
        if (pre[i]!=-1)
            ret+=mat[pre[i]][i];
    return ret:
}
```

四. 图论_网络流

1. 上下界最大流(邻接表形式)

```
for (p=q=0; p <= q\&\&!pre[sink]; t=que[p++])
            for (r=0; r<e[t].size();++r){
                i=e[t][r];
                if (!pre[i]&&(j=mat[t][i]-flow[t][i]))
                    pre[que[q++]=i]=t+1,d[i]=d[t]<j?d[t]:j;</pre>
               else if (!pre[i]\&\&(j=flow[i][t]))
                    pre[que[q++]=i]=-t-1,d[i]=d[t]<j?d[t]:j;</pre>
            }
        if (!pre[sink]) break;
        for (i=sink;i!=source;)
            if (pre[i]>0)
                flow[pre[i]-1][i]+=d[sink],i=pre[i]-1;
            else
                flow[i][-pre[i]-1]-=d[sink],i=-pre[i]-1;
    for (j=i=0;i<n;j+=flow[source][i++]);</pre>
    return j;
}
int limit_max_flow(int n,int mat[][MAXN],int bf[][MAXN],int source,int sink,int
flow[][MAXN]){
    int i,j,sk,ks;
    if (source==sink) return inf;
    for (mat[n][n+1]=mat[n+1][n]=mat[n][n]=mat[n+1][n+1]=i=0;i< n;i++)
        for (mat[n][i]=mat[i][n]=mat[n+1][i]=mat[i][n+1]=j=0; j< n; j++)
            mat[i][j]-=bf[i][j],mat[n][i]+=bf[j][i],mat[i][n+1]+=bf[i][j];
    sk=mat[source][sink],ks=mat[sink][source],mat[source][sink]=mat[sink][source]=inf;
    for (i=0; i< n+2; i++)
        for (j=0;j<n+2;flow[i][j++]=0);
    _max_flow(n+2,mat,n,n+1,flow);
    for (i=0;i< n;i++)
        if (flow[n][i]<mat[n][i]) return -1;</pre>
    flow[source][sink]=flow[sink][source]=0,mat[source][sink]=sk,mat[sink][source]=ks;
    _max_flow(n,mat,source,sink,flow);
    for (i=0;i<n;i++)
       for (j=0; j< n; j++)
            mat[i][j]+=bf[i][j],flow[i][j]+=bf[i][j];
    for (j=i=0;i<n;j+=flow[source][i++]);</pre>
    return j;
}
                             2. 上下界最大流(邻接阵形式)
```

```
//求上下界网络最大流,邻接阵形式 //返回最大流量,-1表示无可行流,flow返回每条边的流量 //传入网络节点数 n,容量 mat,流量下界 bf,源点 source,汇点 sink //MAXN 应比最大结点数多 2,无可行流返回-1 时 mat 未复原!
```

```
#define MAXN 100
#define inf 1000000000
void _max_flow(int n,int mat[][MAXN],int source,int sink,int flow[][MAXN]){
    int pre[MAXN],que[MAXN],d[MAXN],p,q,t,i,j;
    for (;;){
       for (i=0;i<n;pre[i++]=0);
        pre[t=source]=source+1,d[t]=inf;
        for (p=q=0; p <= q \& ! pre[sink]; t = que[p++])
            for (i=0;i<n;i++)
                if (!pre[i]&&j=mat[t][i]-flow[t][i])
                    pre[que[q++]=i]=t+1,d[i]=d[t]<j?d[t]:j;</pre>
                else if (!pre[i]&&j=flow[i][t])
                    pre[que[q++]=i]=-t-1,d[i]=d[t]<j?d[t]:j;</pre>
        if (!pre[sink]) break;
        for (i=sink;i!=source;)
            if (pre[i]>0)
                flow[pre[i]-1][i]+=d[sink],i=pre[i]-1;
            else
                flow[i][-pre[i]-1]-=d[sink],i=-pre[i]-1;
    }
}
int limit_max_flow(int n,int mat[][MAXN],int bf[][MAXN],int source,int sink,int
flow[][MAXN]){
    int i,j,sk,ks;
    if (source==sink) return inf;
    for (mat[n][n+1]=mat[n+1][n]=mat[n][n]=mat[n+1][n+1]=i=0;i< n;i++)
        for (mat[n][i]=mat[i][n]=mat[n+1][i]=mat[i][n+1]=j=0; j< n; j++)
            mat[i][j]-=bf[i][j],mat[n][i]+=bf[j][i],mat[i][n+1]+=bf[i][j];
    sk=mat[source][sink],ks=mat[sink][source],mat[source][sink]=mat[sink][source]=inf;
    for (i=0; i< n+2; i++)
        for (j=0; j< n+2; flow[i][j++]=0);
    max flow(n+2,mat,n,n+1,flow);
    for (i=0;i<n;i++)
        if (flow[n][i]<mat[n][i]) return -1;</pre>
    flow[source][sink]=flow[sink][source]=0,mat[source][sink]=sk,mat[sink][source]=ks;
    _max_flow(n,mat,source,sink,flow);
    for (i=0; i< n; i++)
        for (j=0; j< n; j++)
            mat[i][j]+=bf[i][j],flow[i][j]+=bf[i][j];
    for (j=i=0;i<n;j+=flow[source][i++]);</pre>
    return j;
}
```

3. 上下界最小流(邻接表形式)

```
//求上下界网络最小流,邻接阵形式
//返回最大流量,-1表示无可行流,flow返回每条边的流量
//传入网络节点数 n,容量 mat,流量下界 bf,源点 source,汇点 sink
//MAXN 应比最大结点数多 2, 无可行流返回 - 1 时 mat 未复原!
#define MAXN 100
#define inf 1000000000
int max flow(int n,int mat[][MAXN],int source,int sink,int flow[][MAXN]){
   int pre[MAXN],que[MAXN],d[MAXN],p,q,t,i,j,r;
   vector<int> e[MAXN];
   for (i=0;i< n;i++)
       for (e[i].clear(),j=0;j<n;j++)</pre>
           if (mat[i][j]) e[i].push_back(j),e[j].push_back(i);
   for (;;){
       for (i=0;i<n;pre[i++]=0);
       pre[t=source]=source+1,d[t]=inf;
       for (p=q=0; p <= q \& ! pre[sink]; t = que[p++])
           for (r=0;r<e[t].size();++r){</pre>
               i=e[t][r];
               if (!pre[i]&&(j=mat[t][i]-flow[t][i]))
                   pre[que[q++]=i]=t+1,d[i]=d[t]<j?d[t]:j;</pre>
               else if (!pre[i]&&(j=flow[i][t]))
                   pre[que[q++]=i]=-t-1,d[i]=d[t]<j?d[t]:j;
       if (!pre[sink]) break;
       for (i=sink;i!=source;)
           if (pre[i]>0)
               flow[pre[i]-1][i]+=d[sink],i=pre[i]-1;
           else
               flow[i][-pre[i]-1]-=d[sink],i=-pre[i]-1;
   for (j=i=0;i<n;j+=flow[source][i++]);</pre>
   return j;
}
int limit min flow(int n,int mat[][MAXN],int bf[][MAXN],int source,int sink,int
flow[][MAXN]){
   int i,j,sk,ks;
   if (source==sink) return inf;
   for (mat[n][n+1]=mat[n+1][n]=mat[n][n]=mat[n+1][n+1]=i=0;i< n;i++)
       for (mat[n][i]=mat[i][n]=mat[n+1][i]=mat[i][n+1]=j=0;j< n;j++)
           mat[i][j]-=bf[i][j],mat[n][i]+=bf[j][i],mat[i][n+1]+=bf[i][j];
    sk=mat[source][sink],ks=mat[sink][source],mat[source][sink]=mat[sink][source]=inf;
   for (i=0; i< n+2; i++)
       for (j=0;j<n+2;flow[i][j++]=0);
   _max_flow(n+2,mat,n,n+1,flow);
```

```
if (flow[n][i]<mat[n][i]) return -1;</pre>
   flow[source][sink]=flow[sink][source]=0, mat[source][sink]=sk, mat[sink][source]=ks;
   _max_flow(n,mat,sink,source,flow);
   for (i=0;i< n;i++)
       for (j=0; j< n; j++)
           mat[i][j]+=bf[i][j],flow[i][j]+=bf[i][j];
   for (j=i=0;i<n;j+=flow[source][i++]);</pre>
    return j;
}
                            4. 上下界最小流(邻接阵形式)
//求上下界网络最小流,邻接阵形式
//返回最大流量,-1表示无可行流,flow返回每条边的流量
//传入网络节点数 n,容量 mat,流量下界 bf,源点 source,汇点 sink
//MAXN 应比最大结点数多 2, 无可行流返回-1 时 mat 未复原!
#define MAXN 100
#define inf 1000000000
void max flow(int n,int mat[][MAXN],int source,int sink,int flow[][MAXN]){
   int pre[MAXN],que[MAXN],d[MAXN],p,q,t,i,j;
   for (;;){
       for (i=0;i<n;pre[i++]=0);
       pre[t=source]=source+1,d[t]=inf;
       for (p=q=0; p \le q\&\{pre[sink]; t=que[p++])
           for (i=0;i<n;i++)
               if (!pre[i]&&j=mat[t][i]-flow[t][i])
                   pre[que[q++]=i]=t+1,d[i]=d[t]<j?d[t]:j;</pre>
               else if (!pre[i]&&j=flow[i][t])
                   pre[que[q++]=i]=-t-1,d[i]=d[t]<j?d[t]:j;
       if (!pre[sink]) break;
       for (i=sink;i!=source;)
           if (pre[i]>0)
               flow[pre[i]-1][i]+=d[sink],i=pre[i]-1;
           else
               flow[i][-pre[i]-1]-=d[sink],i=-pre[i]-1;
   }
}
int limit_min_flow(int n,int mat[][MAXN],int bf[][MAXN],int source,int sink,int
flow[][MAXN]){
   int i,j,sk,ks;
   if (source==sink) return inf;
   for (mat[n][n+1]=mat[n+1][n]=mat[n][n]=mat[n+1][n+1]=i=0;i< n;i++)
       for (mat[n][i]=mat[i][n]=mat[n+1][i]=mat[i][n+1]=j=0; j< n; j++)
```

for (i=0;i<n;i++)

```
mat[i][j]-=bf[i][j],mat[n][i]+=bf[j][i],mat[i][n+1]+=bf[i][j];
    sk=mat[source][sink],ks=mat[sink][source],mat[source][sink]=mat[sink][source]=inf;
   for (i=0; i< n+2; i++)
       for (j=0; j< n+2; flow[i][j++]=0);
   _max_flow(n+2,mat,n,n+1,flow);
   for (i=0; i< n; i++)
       if (flow[n][i]<mat[n][i]) return -1;</pre>
   flow[source][sink]=flow[sink][source]=0, mat[source][sink]=sk, mat[sink][source]=ks;
   max flow(n,mat,sink,source,flow);
   for (i=0;i<n;i++)
       for (j=0; j< n; j++)
           mat[i][j]+=bf[i][j],flow[i][j]+=bf[i][j];
   for (j=i=0;i<n;j+=flow[source][i++]);</pre>
   return j;
}
                                5. 最大流(邻接表形式)
//求网络最大流,邻接表形式
//返回最大流量,flow返回每条边的流量
//传入网络节点数 n,容量 mat,邻接表 list,源点 source,汇点 sink
//list[i](vector<int>)存放所有以i相邻的点,包括反向边!!!
#define MAXN 100
#define inf 1000000000
int max flow(int n, int mat[][MAXN], vector<int> list[], int source, int sink, int flow[][MAXN]){
   int pre[MAXN],que[MAXN],d[MAXN],p,q,t,i,j,r;
   if (source==sink) return inf;
   for (i=0;i<n;i++)
       for (j=0; j<n; flow[i][j++]=0);
   for (;;){
       for (i=0;i<n;pre[i++]=0);
       pre[t=source]=source+1,d[t]=inf;
       for (p=q=0; p \le q\&! pre[sink]; t=que[p++])
            for (r=0;r<list[t].size();++r){</pre>
               i=list[t][r];
               if (!pre[i]&&j=mat[t][i]-flow[t][i])
                   pre[que[q++]=i]=t+1,d[i]=d[t]<j?d[t]:j;</pre>
               else if (!pre[i]&&j=flow[i][t])
                   pre[que[q++]=i]=-t-1,d[i]=d[t]<j?d[t]:j;
       if (!pre[sink]) break;
       for (i=sink;i!=source;)
            if (pre[i]>0)
               flow[pre[i]-1][i]+=d[sink],i=pre[i]-1;
           else
```

```
flow[i][-pre[i]-1]-=d[sink],i=-pre[i]-1;
    for (j=i=0;i<n;j+=flow[source][i++]);</pre>
    return j;
}
                         6. 最大流(邻接表形式,邻接阵接口)
//求网络最大流,邻接表形式
//返回最大流量,flow返回每条边的流量
//传入网络节点数 n,容量 mat,源点 source,汇点 sink
#define MAXN 100
#define inf 1000000000
int max_flow(int n,int mat[][MAXN],int source,int sink,int flow[][MAXN]){
    int pre[MAXN],que[MAXN],d[MAXN],p,q,t,i,j,r;
    vector<int> e[MAXN];
    if (source==sink) return inf;
    for (i=0;i<n;i++)
       for (j=0; j< n; flow[i][j++]=0);
    //e[i]存放所有以i相邻的点,包括反向边!!!
    for (i=0;i<n;i++)
       for (e[i].clear(),j=0;j<n;j++)</pre>
           if (mat[i][j]) e[i].push_back(j),e[j].push_back(i);
    for (;;){
       for (i=0;i<n;pre[i++]=0);
       pre[t=source]=source+1,d[t]=inf;
       for (p=q=0; p <= q\&\&! pre[sink]; t=que[p++])
           for (r=0; r<e[t].size();++r){
               i=e[t][r];
               if (!pre[i]&&(j=mat[t][i]-flow[t][i]))
                   pre[que[q++]=i]=t+1,d[i]=d[t]<j?d[t]:j;</pre>
               else if (!pre[i]&&(j=flow[i][t]))
                   pre[que[q++]=i]=-t-1,d[i]=d[t]<j?d[t]:j;
           }
       if (!pre[sink]) break;
       for (i=sink;i!=source;)
           if (pre[i]>0)
               flow[pre[i]-1][i]+=d[sink],i=pre[i]-1;
           else
               flow[i][-pre[i]-1]-=d[sink],i=-pre[i]-1;
    for (j=i=0;i<n;j+=flow[source][i++]);</pre>
    return j;
```

}

7. 最大流(邻接阵形式)

```
//求网络最大流,邻接阵形式
//返回最大流量,flow返回每条边的流量
//传入网络节点数 n,容量 mat,源点 source,汇点 sink
#define MAXN 100
#define inf 1000000000
int max_flow(int n,int mat[][MAXN],int source,int sink,int flow[][MAXN]){
   int pre[MAXN],que[MAXN],d[MAXN],p,q,t,i,j;
   if (source==sink) return inf;
   for (i=0; i< n; i++)
       for (j=0;j<n;flow[i][j++]=0);</pre>
   for (;;){
       for (i=0;i<n;pre[i++]=0);
       pre[t=source]=source+1,d[t]=inf;
       for (p=q=0; p \le q \le ! pre[sink]; t = que[p++])
           for (i=0;i<n;i++)
               if (!pre[i]&&j=mat[t][i]-flow[t][i])
                   pre[que[q++]=i]=t+1,d[i]=d[t]<j?d[t]:j;</pre>
               else if (!pre[i]&&j=flow[i][t])
                   pre[que[q++]=i]=-t-1,d[i]=d[t]<j?d[t]:j;
       if (!pre[sink]) break;
       for (i=sink;i!=source;)
           if (pre[i]>0)
               flow[pre[i]-1][i]+=d[sink],i=pre[i]-1;
           else
               flow[i][-pre[i]-1]-=d[sink],i=-pre[i]-1;
   for (j=i=0;i<n;j+=flow[source][i++]);</pre>
    return j;
}
                            8. 最大流无流量(邻接阵形式)
//求网络最大流,邻接阵形式
//返回最大流量
//传入网络节点数 n,容量 mat,源点 source,汇点 sink
//注意 mat 矩阵被修改
#define MAXN 100
#define inf 1000000000
int max_flow(int n,int mat[][MAXN],int source,int sink){
   int v[MAXN],c[MAXN],p[MAXN],ret=0,i,j;
   for (;;){
```

```
for (i=0;i<n;i++)
           v[i]=c[i]=0;
       for (c[source]=inf;;){
           for (j=-1,i=0;i< n;i++)
               if (!v[i]\&\&c[i]\&\&(j=-1||c[i]>c[j]))
                   j=i;
           if (j<0) return ret;
           if (j==sink) break;
           for (v[j]=1,i=0;i< n;i++)
               if (mat[j][i]>c[i]&&c[j]>c[i])
                   c[i]=mat[j][i]<c[j]?mat[j][i]:c[j],p[i]=j;</pre>
       for (ret+=j=c[i=sink];i!=source;i=p[i])
           mat[p[i]][i]-=j,mat[i][p[i]]+=j;
   }
}
                           9. 最小费用最大流(邻接阵形式)
//求网络最小费用最大流,邻接阵形式
//返回最大流量,flow返回每条边的流量,netcost返回总费用
//传入网络节点数 n,容量 mat,单位费用 cost,源点 source,汇点 sink
#define MAXN 100
#define inf 1000000000
int min cost max flow(int n,int mat[][MAXN],int cost[][MAXN],int source,int sink,int
flow[][MAXN],int& netcost){
   int pre[MAXN],min[MAXN],d[MAXN],i,j,t,tag;
   if (source==sink) return inf;
   for (i=0;i<n;i++)
       for (j=0; j<n; flow[i][j++]=0);
   for (netcost=0;;){
       for (i=0;i<n;i++)
           pre[i]=0,min[i]=inf;
       //通过 bellman ford 寻找最短路,即最小费用可改进路
       for (pre[source]=source+1,min[source]=0,d[source]=inf,tag=1;tag;)
           for (tag=t=0;t<n;t++)
               if (d[t])
                   for (i=0;i<n;i++)
                       if (j=mat[t][i]-flow[t][i]&&min[t]+cost[t][i]<min[i])</pre>
   tag=1,min[i]=min[t]+cost[t][i],pre[i]=t+1,d[i]=d[t]<j?d[t]:j;
                       else if (j=flow[i][t]&&min[t]<inf&&min[t]-cost[i][t]<min[i])</pre>
   tag=1,min[i]=min[t]-cost[i][t],pre[i]=-t-1,d[i]=d[t]<j?d[t]:j;</pre>
       if (!pre[sink]) break;
```

五. 图论_最短路径

1. 最短路径(单源 bellman_ford 邻接阵形式)

```
//单源最短路径,bellman_ford 算法,邻接阵形式,复杂度 0(n^3)
//求出源 s 到所有点的最短路经,传入图的大小 n 和邻接阵 mat
//返回到各点最短距离 min[]和路径 pre[], pre[i]记录 s 到 i 路径上 i 的父结点, pre[s]=-1
//可更改路权类型,路权可为负,若图包含负环则求解失败,返回0
//优化:先删去负边使用 dijkstra 求出上界,加速迭代过程
#define MAXN 200
#define inf 1000000000
typedef int elem t;
int bellman_ford(int n,elem_t mat[][MAXN],int s,elem_t* min,int* pre){
   int v[MAXN],i,j,k,tag;
   for (i=0; i< n; i++)
       min[i]=inf,v[i]=0,pre[i]=-1;
   for (\min[s]=0, j=0; j< n; j++){
       for (k=-1,i=0;i< n;i++)
           if (!v[i]&&(k==-1||min[i]<min[k]))
               k=i:
       for (v[k]=1,i=0;i< n;i++)
           if (!v[i]&&mat[k][i]>=0&&min[k]+mat[k][i]<min[i])</pre>
               min[i]=min[k]+mat[pre[i]=k][i];
   }
   for (tag=1, j=0; tag\&\&j <= n; j++)
       for (tag=i=0;i<n;i++)</pre>
           for (k=0; k< n; k++)
               if (min[k]+mat[k][i]<min[i])</pre>
                   min[i]=min[k]+mat[pre[i]=k][i],tag=1;
    return j<=n;
}
```

2. 最短路径(单源 dijkstra bfs 邻接表形式)

//单源最短路径,用于路权相等的情况,dijkstra 优化为 bfs,邻接表形式,复杂度 0(m)

```
//求出源 s 到所有点的最短路经,传入图的大小 n 和邻接表 list,边权值 len
//返回到各点最短距离 min[]和路径 pre[], pre[i]记录 s 到 i 路径上 i 的父结点, pre[s]=-1
//可更改路权类型,但必须非负且相等!
#define MAXN 200
#define inf 1000000000
typedef int elem t;
struct edge_t{
   int from, to;
   edge_t* next;
};
void dijkstra(int n,edge_t* list[],elem_t len,int s,elem_t* min,int* pre){
   edge t* t;
   int i,que[MAXN],f=0,r=0,p=1,l=1;
   for (i=0;i<n;i++)
       min[i]=inf;
   min[que[0]=s]=0, pre[s]=-1;
   for (;r<=f;l++,r=f+1,f=p-1)
       for (i=r;i<=f;i++)</pre>
           for (t=list[que[i]];t;t=t->next)
               if (min[t->to]==inf)
                  min[que[p++]=t->to]=len*l,pre[t->to]=que[i];
}
                   3. 最短路径(单源 dijkstra bfs 正向表形式)
//单源最短路径,用于路权相等的情况,dijkstra 优化为 bfs,正向表形式,复杂度 O(m)
//求出源 s 到所有点的最短路经,传入图的大小 n 和正向表 list, buf, 边权值 len
//返回到各点最短距离 min[]和路径 pre[], pre[i]记录 s 到 i 路径上 i 的父结点, pre[s]=-1
//可更改路权类型,但必须非负且相等!
#define MAXN 200
#define inf 1000000000
typedef int elem_t;
void dijkstra(int n,int* list,int* buf,elem_t len,int s,elem_t* min,int* pre){
   int i,que[MAXN],f=0,r=0,p=1,l=1,t;
   for (i=0; i< n; i++)
       min[i]=inf;
   min[que[0]=s]=0, pre[s]=-1;
   for (;r \le f;l++,r=f+1,f=p-1)
       for (i=r;i<=f;i++)</pre>
           for (t=list[que[i]];t<list[que[i]+1];t++)</pre>
               if (min[buf[t]]==inf)
                  min[que[p++]=buf[t]]=len*l,pre[buf[t]]=que[i];
}
```

```
//单源最短路径,dijkstra 算法+二分堆,邻接表形式,复杂度 O(mlogm)
//求出源 s 到所有点的最短路经,传入图的大小 n 和邻接表 list
//返回到各点最短距离 min[]和路径 pre[], pre[i]记录 s 到 i 路径上 i 的父结点, pre[s]=-1
//可更改路权类型,但必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem_t;
struct edge_t{
   int from, to;
   elem_t len;
   edge_t* next;
};
#define _{cp(a,b)} ((a).d<(b).d)
struct heap_t{elem_t d;int v;};
struct heap{
   heap_t h[MAXN*MAXN];
   int n,p,c;
   void init(){n=0;}
   void ins(heap_t e){
       for (p=++n;p>1&\&\_cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
       h[p]=e;
   }
   int del(heap_t& e){
       if (!n) return 0;
(e=h[p=1], c=2; c<n&&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
       h[p]=h[n--];return 1;
   }
};
void dijkstra(int n,edge_t* list[],int s,elem_t* min,int* pre){
   heap h;
   edge t* t;heap t e;
   int v[MAXN],i;
   for (i=0;i<n;i++)
       min[i]=inf,v[i]=0,pre[i]=-1;
   h.init();min[e.v=s]=e.d=0,h.ins(e);
   while (h.del(e))
       if (!v[e.v])
           for (v[e.v]=1,t=list[e.v];t;t=t->next)
               if (!v[t->to]&&min[t->from]+t->len<min[t->to])
                   pre[t->to]=t->from,min[e.v=t->to]=e.d=min[t->from]+t->len,h.ins(e);
}
```

5. 最短路径(单源 dijkstra+binary_heap 正向表形式)

```
//单源最短路径,dijkstra 算法+二分堆,正向表形式,复杂度 0(mlogm)
//求出源 s 到所有点的最短路经,传入图的大小 n 和正向表 list, buf
//返回到各点最短距离 min[]和路径 pre[], pre[i]记录 s 到 i 路径上 i 的父结点, pre[s]=-1
//可更改路权类型,但必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem_t;
struct edge_t{
    int to;
    elem_t len;
};
#define cp(a,b) ((a).d<(b).d)
struct heap_t{elem_t d;int v;};
struct heap{
    heap_t h[MAXN*MAXN];
    int n,p,c;
   void init(){n=0;}
   void ins(heap_t e){
       for (p=++n;p>1&\&_cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
       h[p]=e;
    }
    int del(heap_t& e){
       if (!n) return 0;
(e=h[p=1], c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1], h[c]))], h[n]); h[p]=h[c], p=c, c<<=1);
       h[p]=h[n--];return 1;
    }
};
void dijkstra(int n,int* list,edge_t* buf,int s,elem_t* min,int* pre){
    heap h;heap_t e;
    int v[MAXN],i,t,f;
    for (i=0; i< n; i++)
       min[i]=inf,v[i]=0,pre[i]=-1;
    h.init(); min[e.v=s]=e.d=0,h.ins(e);
   while (h.del(e))
       if (!v[e.v])
           for (v[f=e.v]=1,t=list[f];t<list[f+1];t++)</pre>
               if (!v[buf[t].to]&&min[f]+buf[t].len<min[buf[t].to])</pre>
                   pre[buf[t].to]=f,min[e.v=buf[t].to]=e.d=min[f]+buf[t].len,h.ins(e);
}
```

6. 最短路径(单源 dijkstra+mapped heap 邻接表形式)

//单源最短路径,dijkstra 算法+映射二分堆,邻接表形式,复杂度 0(mlogn) //求出源 s 到所有点的最短路经,传入图的大小 n 和邻接表 list

```
//返回到各点最短距离 min[]和路径 pre[], pre[i]记录 s 到 i 路径上 i 的父结点, pre[s]=-1
//可更改路权类型,但必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem_t;
struct edge_t{
    int from, to;
    elem_t len;
   edge_t* next;
};
#define _{cp(a,b)} ((a)<(b))
struct heap{
    elem_t h[MAXN+1];
    int ind[MAXN+1],map[MAXN+1],n,p,c;
    void init(){n=0;}
   void ins(int i,elem_t e){
        for (p=++n;p>1&\&\_cp(e,h[p>>1]);h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
        h[map[ind[p]=i]=p]=e;
    }
    int del(int i,elem_t& e){
        i=map[i];if (i<1||i>n) return 0;
        for (e=h[p=i];p>1;h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
        for
(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<
=1);
        h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
    }
    int delmin(int& i,elem t& e){
        if (n<1) return 0;i=ind[1];</pre>
(e=h[p=1], c=2; c<n&_cp(h[c+=(c<n-1&&_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c]
],p=c,c<<=1);
        h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
    }
};
void dijkstra(int n,edge_t* list[],int s,elem_t* min,int* pre){
    heap h;
    edge_t* t;elem_t e;
    int v[MAXN],i;
    for (h.init(),i=0;i<n;i++)
        min[i]=((i==s)?0:inf), v[i]=0, pre[i]=-1, h.ins(i, min[i]);
   while (h.delmin(i,e))
        for (v[i]=1,t=list[i];t;t=t->next)
            if (!v[t->to]&&min[i]+t->len<min[t->to])
```

```
pre[t->to]=i,h.del(t->to,e),min[t->to]=e=min[i]+t->len,h.ins(t->to,e);
}
```

7. 最短路径(单源 dijkstra+mapped heap 正向表形式)

```
//单源最短路径,dijkstra 算法+映射二分堆,正向表形式,复杂度 0(mlogn)
//求出源 s 到所有点的最短路经,传入图的大小 n 和正向表 list,buf
//返回到各点最短距离 min[]和路径 pre[], pre[i]记录 s 到 i 路径上 i 的父结点, pre[s]=-1
//可更改路权类型,但必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem_t;
struct edge_t{
           int to;
           elem t len;
};
#define _{cp(a,b)}((a)<(b))
struct heap{
           elem_t h[MAXN+1];
           int ind[MAXN+1],map[MAXN+1],n,p,c;
          void init(){n=0;}
           void ins(int i,elem_t e){
                      for (p=++n;p>1&\&\_cp(e,h[p>>1]);h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
                      h[map[ind[p]=i]=p]=e;
           }
           int del(int i,elem t& e){
                      i=map[i];if (i<1||i>n) return 0;
                      for (e=h[p=i];p>1;h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
                      for
(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[n]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[
=1);
                      h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
           int delmin(int& i,elem_t& e){
                      if (n<1) return 0;i=ind[1];</pre>
(e=h[p=1],c=2;c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c]
],p=c,c<<=1);
                      h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
           }
};
void dijkstra(int n,int* list,edge_t* buf,int s,elem_t* min,int* pre){
           heap h;elem_t e;
           int v[MAXN],i,t;
           for (h.init(),i=0;i<n;i++)
```

```
min[i]=((i==s)?0:inf),v[i]=0,pre[i]=-1,h.ins(i,min[i]);
         while (h.delmin(i,e))
                   for (v[i]=1,t=list[i];t<list[i+1];t++)</pre>
                            if (!v[buf[t].to]&&min[i]+buf[t].len<min[buf[t].to])</pre>
         pre[buf[t].to]=i,h.del(buf[t].to,e),min[buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t].len,h.ins(buf[t].to]=e=min[i]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+buf[t]+
to,e);
}
                                                       8. 最短路径(单源 dijkstra 邻接阵形式)
//单源最短路径,dijkstra 算法,邻接阵形式,复杂度 0(n^2)
//求出源 s 到所有点的最短路经,传入图的顶点数 n, (有向)邻接矩阵 mat
//返回到各点最短距离 min[]和路径 pre[],pre[i]记录 s 到 i 路径上 i 的父结点,pre[s]=-1
//可更改路权类型,但必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem_t;
void dijkstra(int n,elem_t mat[][MAXN],int s,elem_t* min,int* pre){
         int v[MAXN],i,j,k;
         for (i=0; i< n; i++)
                  min[i]=inf,v[i]=0,pre[i]=-1;
         for (\min[s]=0, j=0; j< n; j++){
                   for (k=-1, i=0; i< n; i++)
                            if (!v[i]&&(k==-1||min[i]<min[k]))
                                     k=i;
                   for (v[k]=1,i=0;i< n;i++)
                            if (!v[i]&&min[k]+mat[k][i]<min[i])</pre>
                                     min[i]=min[k]+mat[pre[i]=k][i];
         }
}
                                             9. 最短路径(多源 floyd warshall 邻接阵形式)
//多源最短路径,floyd warshall 算法,复杂度 0(n^3)
//求出所有点对之间的最短路经,传入图的大小和邻接阵
//返回各点间最短距离 min[]和路径 pre[], pre[i][j]记录 i 到 j 最短路径上 j 的父结点
//可更改路权类型,路权必须非负!
#define MAXN 200
#define inf 1000000000
typedef int elem_t;
void floyd_warshall(int n,elem_t mat[][MAXN],elem_t min[][MAXN],int pre[][MAXN]){
         int i,j,k;
         for (i=0;i<n;i++)
                   for (j=0; j< n; j++)
                            min[i][j]=mat[i][j],pre[i][j]=(i==j)?-1:i;
```

```
for (k=0; k< n; k++)
        for (i=0;i<n;i++)
             for (j=0; j< n; j++)
                 if (min[i][k]+min[k][j]<min[i][j])</pre>
                     min[i][j]=min[i][k]+min[k][j],pre[i][j]=pre[k][j];
}
```

六.图论_连通性

1. 无向图关键边(dfs 邻接阵形式)

```
//无向图的关键边,dfs邻接阵形式,0(n^2)
//返回关键边条数,key[][2]返回边集
//传入图的大小 n 和邻接阵 mat, 不相邻点边权 0
#define MAXN 100
void search(int n,int mat[][MAXN],int* dfn,int* low,int now,int& cnt,int key[][2]){
    for (low[now]=dfn[now],i=0;i<n;i++)</pre>
       if (mat[now][i]){
           if (!dfn[i]){
               dfn[i]=dfn[now]+1;
               search(n,mat,dfn,low,i,cnt,key);
               if (low[i]>dfn[now])
                   key[cnt][0]=i,key[cnt++][1]=now;
               if (low[i]<low[now])</pre>
                   low[now]=low[i];
           else if (dfn[i]<dfn[now]-1&&dfn[i]<low[now])</pre>
               low[now]=lev[i];
       }
}
int key_edge(int n,int mat[][MAXN],int key[][2]){
    int ret=0,i,dfn[MAXN],low[MAXN];
    for (i=0; i< n; dfn[i++]=0);
    for (i=0;i<n;i++)
       if (!dfn[i])
           dfn[i]=1,bridge(n,mat,dfn,low,i,ret,key);
    return ret;
}
                          2. 无向图关键点(dfs 邻接阵形式)
```

//无向图的关键点,dfs 邻接阵形式,0(n^2) //返回关键点个数,key[]返回点集

```
//传入图的大小 n 和邻接阵 mat, 不相邻点边权 0
#define MAXN 110
void search(int n,int mat[][MAXN],int* dfn,int* low,int now,int& ret,int* key,int& cnt,int
root,int& rd,int* bb){
    int i:
    dfn[now]=low[now]=++cnt;
    for (i=0; i< n; i++)
       if (mat[now][i]){
           if (!dfn[i]){
               search(n,mat,dfn,low,i,ret,key,cnt,root,rd,bb);
               if (low[i]<low[now])</pre>
                   low[now]=low[i];
               if (low[i]>=dfn[now]){
                   if (now!=root&&!bb[now])
                       key[ret++]=now,bb[now]=1;
                   else if(now==root)
                       rd++;
               }
           }
           else if (dfn[i]<low[now])</pre>
               low[now]=dfn[i];
       }
}
int key_vertex(int n,int mat[][MAXN],int* key){
    int ret=0,i,cnt,rd,dfn[MAXN],low[MAXN],bb[MAXN];
    for (i=0;i<n;dfn[i++]=bb[i]=0);
    for (cnt=i=0;i<n;i++)</pre>
       if (!dfn[i]){
            rd=0;
           search(n,mat,dfn,low,i,ret,key,cnt,i,rd,bb);
           if (rd>1&&!bb[i])
               key[ret++]=i,bb[i]=1;
       }
    return ret;
}
                            3. 无向图块(bfs 邻接阵形式)
//无向图的块,dfs 邻接阵形式,0(n^2)
//每产生一个块调用 dummy
//传入图的大小 n 和邻接阵 mat, 不相邻点边权 0
#define MAXN 100
#include <iostream.h>
void dummy(int n,int* a){
```

```
for (int i=0; i< n; i++)
        cout<<a[i]<<' ';
    cout<<endl;</pre>
}
void search(int n,int mat[][MAXN],int* dfn,int* low,int now,int& cnt,int* st,int& sp){
    int i,m,a[MAXN];
    dfn[st[sp++]=now]=low[now]=++cnt;
    for (i=0;i<n;i++)
        if (mat[now][i]){
            if (!dfn[i]){
               search(n,mat,dfn,low,i,cnt,st,sp);
               if (low[i]<low[now])</pre>
                   low[now]=low[i];
               if (low[i]>=dfn[now]){
                   for (st[sp]=-1,a[0]=now,m=1;st[sp]!=i;a[m++]=st[--sp]);
                   dummy(m,a);
               }
            }
            else if (dfn[i]<low[now])</pre>
               low[now]=dfn[i];
       }
}
void block(int n,int mat[][MAXN]){
    int i,cnt,dfn[MAXN],low[MAXN],st[MAXN],sp=0;
    for (i=0; i< n; dfn[i++]=0);
    for (cnt=i=0;i<n;i++)</pre>
        if (!dfn[i])
            search(n,mat,dfn,low,i,cnt,st,sp);
}
                         4. 无向图连通分支(bfs 邻接阵形式)
//无向图连通分支,bfs 邻接阵形式,0(n^2)
//返回分支数,id返回1..分支数的值
//传入图的大小 n 和邻接阵 mat, 不相邻点边权 0
#define MAXN 100
int find_components(int n,int mat[][MAXN],int* id){
  int ret,k,i,j,m;
  for (k=0; k< n; id[k++]=0);
  for (ret=k=0; k< n; k++)
      if (!id[k])
          for (id[k]=-1,ret++,m=1;m;)
              for (m=i=0;i< n;i++)
                  if (id[i] == -1)
```

```
for (m++,id[i]=ret,j=0;j<n;j++)</pre>
                         if (!id[j]&&mat[i][j])
                            id[j]=-1;
  return ret;
}
                        5. 无向图连通分支(dfs 邻接阵形式)
//无向图连通分支,dfs 邻接阵形式,0(n^2)
//返回分支数,id返回1..分支数的值
//传入图的大小 n 和邻接阵 mat, 不相邻点边权 0
#define MAXN 100
void floodfill(int n,int mat[][MAXN],int* id,int now,int tag){
   for (id[now]=tag,i=0;i<n;i++)</pre>
       if (!id[i]&&mat[now][i])
           floodfill(n,mat,id,i,tag);
}
int find_components(int n,int mat[][MAXN],int* id){
   int ret,i;
   for (i=0; i< n; id[i++]=0);
   for (ret=i=0;i<n;i++)</pre>
       if (!id[i])
           floodfill(n,mat,id,i,++ret);
   return ret;
}
                       6. 有向图强连通分支(bfs 邻接阵形式)
//有向图强连通分支,bfs 邻接阵形式,0(n^2)
//返回分支数,id返回1..分支数的值
//传入图的大小 n 和邻接阵 mat, 不相邻点边权 0
#define MAXN 100
int find_components(int n,int mat[][MAXN],int* id){
   int ret=0,a[MAXN],b[MAXN],c[MAXN],d[MAXN],i,j,k,t;
   for (k=0; k< n; id[k++]=0);
   for (k=0; k< n; k++)
       if (!id[k]){
           for (i=0;i<n;i++)
               a[i]=b[i]=c[i]=d[i]=0;
           a[k]=b[k]=1;
           for (t=1;t;)
               for (t=i=0; i< n; i++){
                   if (a[i]&&!c[i])
                      for (c[i]=t=1, j=0; j< n; j++)
```

```
if (mat[i][j]&&!a[j])
                               a[j]=1;
                   if (b[i]&&!d[i])
                       for (d[i]=t=1, j=0; j< n; j++)
                           if (mat[j][i]&&!b[j])
                               b[j]=1;
               }
           for (ret++,i=0;i<n;i++)
               if (a[i]&b[i])
                   id[i]=ret;
    return ret;
}
                        7. 有向图强连通分支(dfs 邻接阵形式)
//有向图强连通分支,dfs 邻接阵形式,0(n^2)
//返回分支数,id返回1..分支数的值
//传入图的大小 n 和邻接阵 mat, 不相邻点边权 0
#define MAXN 100
void search(int n,int mat[][MAXN],int* dfn,int* low,int now,int& cnt,int& tag,int* id,int*
st,int& sp){
    int i,j;
    dfn[st[sp++]=now]=low[now]=++cnt;
    for (i=0;i<n;i++)
        if (mat[now][i]){
           if (!dfn[i]){
               ssearch(n,mat,dfn,low,i,cnt,tag,id,st,sp);
               if (low[i]<low[now])</pre>
                   low[now]=low[i];
           }
           else if (dfn[i]<dfn[now]){</pre>
               for (j=0;j<sp&&st[j]!=i;j++);
               if (j<cnt&dfn[i]<low[now])</pre>
                   low[now]=dfn[i];
           }
    if (low[now]==dfn[now])
        for (tag++;st[sp]!=now;id[st[--sp]]=tag);
}
int find components(int n,int mat[][MAXN],int* id){
    int ret=0,i,cnt,sp,st[MAXN],dfn[MAXN],low[MAXN];
    for (i=0; i< n; dfn[i++]=0);
    for (sp=cnt=i=0;i<n;i++)</pre>
        if (!dfn[i])
```

```
search(n,mat,dfn,low,i,cnt,ret,id,st,sp);
   return ret;
}
                         8. 有向图最小点基(邻接阵形式)
//有向图最小点基,邻接阵形式,0(n^2)
//返回电集大小和点集
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
//需要调用强连通分支
#define MAXN 100
int base_vertex(int n,int mat[][MAXN],int* sets){
   int ret=0,id[MAXN],v[MAXN],i,j;
   j=find components(n,mat,id);
   for (i=0; i< j; v[i++]=1);
   for (i=0;i<n;i++)</pre>
       for (j=0; j< n; j++)
           if (id[i]!=id[j]&&mat[i][j])
```

v[id[j]-1]=0;

v[id[sets[ret++]=i]-1]=0;

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return ret;

}

for (i=0;i<n;i++)

if (v[id[i]-1])

1. 欧拉回路(邻接阵形式)

```
//求欧拉回路或欧拉路,邻接阵形式,复杂度 0(n^2)
//返回路径长度,path 返回路径(有向图时得到的是反向路径)
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
//可以有自环与重边,分为无向图和有向图

#define MAXN 100

void find_path_u(int n,int mat[][MAXN],int now,int& step,int* path){
    int i;
    for (i=n-1;i>=0;i--)
        while (mat[now][i]){
            mat[now][i]--,mat[i][now]--;
            find_path_u(n,mat,i,step,path);
        }
    path[step++]=now;
}
```

```
void find_path_d(int n,int mat[][MAXN],int now,int& step,int* path){
   int i;
   for (i=n-1; i>=0; i--)
       while (mat[now][i]){
           mat[now][i]--;
           find_path_d(n,mat,i,step,path);
       }
   path[step++]=now;
}
int euclid_path(int n,int mat[][MAXN],int start,int* path){
   int ret=0;
   find_path_u(n,mat,start,ret,path);
// find_path_d(n,mat,start,ret,path);
   return ret;
}
                                   2. 前序表转化
//将用边表示的树转化为前序表示的树
//传入节点数 n 和邻接表 list[],邻接表必须是双向的,会在函数中释放
//pre[]返回前序表,map[]返回前序表中的节点到原来节点的映射
#define MAXN 10000
struct node{
   int to;
   node* next;
};
void prenode(int n,node* list[],int* pre,int* map,int* v,int now,int last,int& id){
   node* t;
   int p=id++;
   for (v[map[p]=now]=1,pre[p]=last;list[now];){
       t=list[now],list[now]=t->next;
       if (!v[t->to])
           prenode(n,list,pre,map,v,t->to,p,id);
   }
}
void makepre(int n,node* list[],int* pre,int* map){
   int v[MAXN],id=0,i;
   for (i=0; i< n; v[i++]=0);
   prenode(n,list,pre,map,v,0,-1,id);
}
```

3. 树的优化算法

```
int max_node_independent(int n,int* pre,int* set){
    int c[MAXN],i,ret=0;
    for (i=0;i<n;i++)</pre>
        c[i]=set[i]=0;
   for (i=n-1;i>=0;i--)
        if (!c[i]){
           set[i]=1;
            if (pre[i]!=-1)
               c[pre[i]]=1;
            ret++;
    return ret;
}
//最大边独立集
int max_edge_independent(int n,int* pre,int* set){
    int c[MAXN],i,ret=0;
    for (i=0;i<n;i++)
        c[i]=set[i]=0;
    for (i=n-1; i>=0; i--)
        if (!c[i]&&pre[i]!=-1&&!c[pre[i]]){
            set[i]=1;
           c[pre[i]]=1;
            ret++;
       }
    return ret;
}
//最小顶点覆盖集
int min_node_cover(int n,int* pre,int* set){
    int c[MAXN],i,ret=0;
    for (i=0;i<n;i++)
        c[i]=set[i]=0;
    for (i=n-1;i>=0;i--)
        if (!c[i]&&pre[i]!=-1&&!c[pre[i]]){
            set[i]=1;
           c[pre[i]]=1;
            ret++;
        }
    return ret;
}
//最小顶点支配集
int min_node_dominant(int n,int* pre,int* set){
    int c[MAXN],i,ret=0;
    for (i=0;i<n;i++)</pre>
```

```
c[i]=set[i]=0;
   for (i=n-1;i>=0;i--)
       if (!c[i]&&(pre[i]==-1||!set[pre[i]])){
           if (pre[i]!=-1){
               set[pre[i]]=1;
               c[pre[i]]=1;
               if (pre[pre[i]]!=-1)
                  c[pre[pre[i]]]=1;
           }
           else
               set[i]=1;
           ret++;
       }
   return ret;
}
                             4. 拓扑排序(邻接阵形式).
//拓扑排序,邻接阵形式,复杂度 0(n^2)
//如果无法完成排序,返回0,否则返回1,ret 返回有序点列
//传入图的大小 n 和邻接阵 mat,不相邻点边权 0
#define MAXN 100
int toposort(int n,int mat[][MAXN],int* ret){
   int d[MAXN],i,j,k;
   for (i=0;i<n;i++)
       for (d[i]=j=0; j< n; d[i]+=mat[j++][i]);
   for (k=0; k< n; ret[k++]=i){
       for (i=0;d[i]\&\&i<n;i++);
       if (i==n)
           return 0;
       for (d[i]=-1,j=0;j< n;j++)
           d[j]-=mat[i][j];
   }
   return 1;
}
                                   5. 最佳边割集
//最佳边割集
#define MAXN 100
#define inf 1000000000
int max_flow(int n,int mat[][MAXN],int source,int sink){
   int v[MAXN],c[MAXN],p[MAXN],ret=0,i,j;
   for (;;){
       for (i=0;i<n;i++)
           v[i]=c[i]=0;
```

```
for (c[source]=inf;;){
            for (j=-1, i=0; i< n; i++)
                if (!v[i]\&\&c[i]\&\&(j=-1||c[i]>c[j]))
            if (j<0) return ret;
            if (j==sink) break;
            for (v[j]=1,i=0;i< n;i++)
                if (mat[j][i]>c[i]&&c[j]>c[i])
                     c[i]=mat[j][i]<c[j]?mat[j][i]:c[j],p[i]=j;</pre>
        }
        for (ret+=j=c[i=sink];i!=source;i=p[i])
            mat[p[i]][i]-=j, mat[i][p[i]]+=j;
    }
}
int best_edge_cut(int n,int mat[][MAXN],int source,int sink,int set[][2],int& mincost){
    int m0[MAXN][MAXN],m[MAXN][MAXN],i,j,k,l,ret=0,last;
    if (source==sink)
        return -1;
    for (i=0;i<n;i++)
        for (j=0; j< n; j++)
            m0[i][j]=mat[i][j];
    for (i=0;i< n;i++)
        for (j=0; j< n; j++)
            m[i][j]=m\Theta[i][j];
    mincost=last=max_flow(n,m,source,sink);
    for (k=0; k<n&&last; k++)</pre>
        for (l=0;l<n&&last;l++)</pre>
            if (m0[k][l]){
                for (i=0;i<n+n;i++)
                     for (j=0; j< n+n; j++)
                         m[i][j]=m0[i][j];
                m[k][l]=0;
                if (max flow(n,m,source,sink)==last-mat[k][l]){
                     set[ret][0]=k;
                     set[ret++][1]=l;
                    m0[k][l]=0;
                     last-=mat[k][l];
                }
            }
    return ret;
}
```

6. 最佳顶点割集

//最佳顶点割集 #define MAXN 100

```
int max flow(int n,int mat[][MAXN],int source,int sink){
    int v[MAXN],c[MAXN],p[MAXN],ret=0,i,j;
    for (;;){
        for (i=0;i<n;i++)
            v[i]=c[i]=0;
        for (c[source]=inf;;){
            for (j=-1, i=0; i< n; i++)
                if (!v[i]\&\&c[i]\&\&(j=-1||c[i]>c[j]))
                    j=i;
            if (j<0) return ret;
            if (j==sink) break;
            for (v[j]=1,i=0;i< n;i++)
                if (mat[j][i]>c[i]&&c[j]>c[i])
                    c[i]=mat[j][i]<c[j]?mat[j][i]:c[j],p[i]=j;</pre>
        }
        for (ret+=j=c[i=sink];i!=source;i=p[i])
            mat[p[i]][i]-=j,mat[i][p[i]]+=j;
    }
}
int best_vertex_cut(int n,int mat[][MAXN],int* cost,int source,int sink,int* set,int&
mincost){
    int m0[MAXN][MAXN],m[MAXN][MAXN],i,j,k,ret=0,last;
    if (source==sink||mat[source][sink])
        return -1;
    for (i=0;i<n+n;i++)
        for (j=0; j< n+n; j++)
            m0[i][j]=0;
    for (i=0;i<n;i++)</pre>
        for (j=0; j< n; j++)
            if (mat[i][j])
                m0[i][n+j]=inf;
    for (i=0;i<n;i++)
        m0[n+i][i]=cost[i];
    for (i=0;i<n+n;i++)</pre>
        for (j=0; j< n+n; j++)
            m[i][j]=m0[i][j];
    mincost=last=max_flow(n+n,m,source,n+sink);
    for (k=0; k<n&&last; k++)</pre>
        if (k!=source&&k!=sink){
            for (i=0;i<n+n;i++)
                for (j=0; j< n+n; j++)
                    m[i][j]=mΘ[i][j];
            m[n+k][k]=0;
```

```
if (max_flow(n+n,m,source,n+sink)==last-cost[k]){
                set[ret++]=k;
                m0[n+k][k]=0;
                last-=cost[k];
            }
    return ret;
}
                                      7. 最小边割集
//最小边割集
#define MAXN 100
#define inf 1000000000
int max_flow(int n,int mat[][MAXN],int source,int sink){
    int v[MAXN],c[MAXN],p[MAXN],ret=0,i,j;
    for (;;){
        for (i=0;i< n;i++)
            v[i]=c[i]=0;
        for (c[source]=inf;;){
            for (j=-1, i=0; i< n; i++)
                if (!v[i]\&\&c[i]\&\&(j==-1||c[i]>c[j]))
                    j=i;
            if (j<0) return ret;
            if (j==sink) break;
            for (v[j]=1,i=0;i< n;i++)
                if (mat[j][i]>c[i]\&\&c[j]>c[i])
                    c[i]=mat[j][i]<c[j]?mat[j][i]:c[j],p[i]=j;</pre>
        for (ret+=j=c[i=sink];i!=source;i=p[i])
            mat[p[i]][i]-=j,mat[i][p[i]]+=j;
    }
}
int min edge cut(int n,int mat[][MAXN],int source,int sink,int set[][2]){
    int m0[MAXN][MAXN],m[MAXN][MAXN],i,j,k,l,ret=0,last;
    if (source==sink)
        return -1;
    for (i=0;i<n;i++)
        for (j=0; j< n; j++)
            m0[i][j]=(mat[i][j]!=0);
    for (i=0;i<n;i++)</pre>
        for (j=0; j< n; j++)
            m[i][j]=m\Theta[i][j];
    last=max_flow(n,m,source,sink);
    for (k=0; k<n&&last; k++)</pre>
```

```
for (l=0;l<n&&last;l++)</pre>
            if (m0[k][l]){
                for (i=0;i<n+n;i++)</pre>
                    for (j=0; j< n+n; j++)
                        m[i][j]=m0[i][j];
                m[k][l]=0;
                if (max_flow(n,m,source,sink)<last){</pre>
                    set[ret][0]=k;
                    set[ret++][1]=l;
                    m0[k][l]=0;
                    last--;
                }
            }
    return ret;
}
                                     8. 最小顶点割集
//最小顶点割集
#define MAXN 100
#define inf 1000000000
int max_flow(int n,int mat[][MAXN],int source,int sink){
    int v[MAXN],c[MAXN],p[MAXN],ret=0,i,j;
    for (;;){
        for (i=0;i<n;i++)
            v[i]=c[i]=0;
        for (c[source]=inf;;){
            for (j=-1, i=0; i< n; i++)
                if (!v[i]\&\&c[i]\&\&(j==-1||c[i]>c[j]))
            if (j<0) return ret;
            if (j==sink) break;
            for (v[j]=1,i=0;i< n;i++)
                if (mat[j][i]>c[i]&&c[j]>c[i])
                    c[i]=mat[j][i]<c[j]?mat[j][i]:c[j],p[i]=j;</pre>
        for (ret+=j=c[i=sink];i!=source;i=p[i])
            mat[p[i]][i]-=j,mat[i][p[i]]+=j;
    }
}
int min vertex cut(int n,int mat[][MAXN],int source,int sink,int* set){
    int m0[MAXN][MAXN],m[MAXN][MAXN],i,j,k,ret=0,last;
    if (source==sink||mat[source][sink])
        return -1;
    for (i=0;i<n+n;i++)
```

```
for (j=0; j< n+n; j++)
           m0[i][j]=0;
    for (i=0;i<n;i++)</pre>
       for (j=0; j< n; j++)
           if (mat[i][j])
               m0[i][n+j]=inf;
    for (i=0;i<n;i++)
       m0[n+i][i]=1;
    for (i=0;i<n+n;i++)
       for (j=0; j< n+n; j++)
           m[i][j]=m0[i][j];
    last=max_flow(n+n,m,source,n+sink);
    for (k=0; k<n&&last; k++)</pre>
       if (k!=source&&k!=sink){
            for (i=0; i< n+n; i++)
               for (j=0; j< n+n; j++)
                   m[i][j]=m0[i][j];
           m[n+k][k]=0;
           if (max_flow(n+n,m,source,n+sink)<last){</pre>
               set[ret++]=k;
               m0[n+k][k]=0;
               last--:
           }
    return ret;
}
                                    9. 最小路径覆盖
//最小路径覆盖,0(n^3)
//求解最小的路径覆盖图中所有点,有向图无向图均适用
//注意此问题等价二分图最大匹配,可以用邻接表或正向表减小复杂度
//返回最小路径条数,pre 返回前指针(起点-1),next 返回后指针(终点-1)
#include <string.h>
#define MAXN 310
#define clr(x) memset(x,0xff,sizeof(int)*n)
int hungary(int n,int mat[][MAXN],int* match1,int* match2){
    int s[MAXN],t[MAXN],p,q,ret=0,i,j,k;
    for (_clr(match1),_clr(match2),i=0;i<n;ret+=(match1[i++]>=0))
       for (clr(t), s[p=q=0]=i; p <= q\&match1[i] < 0; p++)
           for (k=s[p],j=0;j<n&&match1[i]<0;j++)</pre>
               if (mat[k][j]\&\&t[j]<0){
                   s[++q]=match2[j],t[j]=k;
                   if (s[q]<0)
                       for (p=j;p>=0;j=p)
                           match2[j]=k=t[j],p=match1[k],match1[k]=j;
```

```
return ret;
}
inline int path_cover(int n,int mat[][MAXN],int* pre,int* next){
   return n-hungary(n,mat,next,pre);
}
```

八. 图论_NP 搜索

1. 最大团(n 小于 64)(faster)

```
/**
* WishingBone's ACM/ICPC Routine Library
* maximum clique solver
#include <vector>
using std::vector;
// clique solver calculates both size and consitution of maximum clique
// uses bit operation to accelerate searching
// graph size limit is 63, the graph should be undirected
// can optimize to calculate on each component, and sort on vertex degrees
// can be used to solve maximum independent set
class clique {
public:
   static const long long ONE = 1;
   static const long long MASK = (1 << 21) - 1;
   char* bits;
   int n, size, cmax[63];
   long long mask[63], cons;
   // initiate lookup table
   clique() {
      bits = new char[1 << 21];
      bits[0] = 0;
      for (int i = 1; i < 1 \iff 21; ++i) bits[i] = bits[i >> 1] + (i & 1);
   ~clique() {
      delete bits;
   // search routine
   bool search(int step, int size, long long more, long long con);
```

```
// solve maximum clique and return size
   int sizeClique(vector<vector<int> >& mat);
   // solve maximum clique and return constitution
   vector<int> consClique(vector<vector<int> >& mat);
};
// search routine
// step is node id, size is current solution, more is available mask, cons is
constitution mask
bool clique::search(int step, int size, long long more, long long cons) {
   if (step >= n) {
      // a new solution reached
      this->size = size;
      this->cons = cons;
      return true;
   long long now = ONE << step;</pre>
   if ((now \& more) > 0) {
      long long next = more & mask[step];
      if (size + bits[next & MASK] + bits[(next >> 21) & MASK] + bits[next >>
42] >= this->size
             && size + cmax[step] > this->size) {
          // the current node is in the clique
          if (search(step + 1, size + 1, next, cons | now)) return true;
      }
   long long next = more & ~now;
   if (size + bits[next & MASK] + bits[(next >> 21) & MASK] + bits[next >> 42]
> this->size) {
      // the current node is not in the clique
      if (search(step + 1, size, next, cons)) return true;
   return false;
}
// solve maximum clique and return size
int clique::sizeClique(vector<vector<int> >& mat) {
   n = mat.size();
   // generate mask vectors
   for (int i = 0; i < n; ++i) {
      mask[i] = 0;
      for (int j = 0; j < n; ++j) if (mat[i][j] > 0) mask[i] |= ONE << j;
   size = 0;
   for (int i = n - 1; i \ge 0; --i) {
      search(i + 1, 1, mask[i], ONE << i);</pre>
```

```
cmax[i] = size;
   }
   return size;
}
// solve maximum clique and return constitution
// calls sizeClique and restore cons
vector<int> clique::consClique(vector<vector<int> >& mat) {
   sizeClique(mat);
   vector<int> ret;
   for (int i = 0; i < n; ++i) if ((cons & (ONE << i)) > 0) ret.push_back(i);
   return ret;
}
                                       2. 最大团
//最大团
//返回最大团大小和一个方案,传入图的大小 n 和邻接阵 mat
//mat[i][j]为布尔量
#define MAXN 60
void clique(int n, int* u, int mat[][MAXN], int size, int& max, int& bb, int* res, int* rr,
int* c) {
   int i, j, vn, v[MAXN];
    if (n) {
       if (size + c[u[0]] \le max) return;
       for (i = 0; i < n + size - max && i < n; ++ i) {
           for (j = i + 1, vn = 0; j < n; ++ j)
               if (mat[u[i]][u[j]])
                   v[vn ++] = u[j];
           rr[size] = u[i];
           clique(vn, v, mat, size + 1, max, bb, res, rr, c);
           if (bb) return;
       }
    } else if (size > max) {
       max = size;
       for (i = 0; i < size; ++ i)
           res[i] = rr[i];
       bb = 1;
   }
}
int maxclique(int n, int mat[][MAXN], int *ret) {
    int max = 0, bb, c[MAXN], i, j;
    int vn, v[MAXN], rr[MAXN];
    for (c[i = n - 1] = 0; i >= 0; -- i) {
```

九.组合

1. 排列组合生成

```
//gen_perm产生字典序排列 P(n,m)
//gen_comb 产生字典序组合 C(n,m)
//gen_perm_swap产生相邻位对换全排列 P(n,n)
//产生元素用 1..n 表示
//dummy 为产生后调用的函数,传入 a[]和 n,a[0]..a[n-1]为一次产生的结果
#define MAXN 100
int count;
#include <iostream.h>
void dummy(int* a,int n){
   int i;
   cout<<count++<<": ";</pre>
   for (i=0;i<n-1;i++)
       cout<<a[i]<<' ';
   cout<<a[n-1]<<endl;</pre>
}
void _gen_perm(int* a,int n,int m,int l,int* temp,int* tag){
   int i;
   if (l==m)
       dummy(temp,m);
   else
       for (i=0;i<n;i++)
           if (!tag[i]){
               temp[l]=a[i],tag[i]=1;
               _gen_perm(a,n,m,l+1,temp,tag);
               tag[i]=0;
           }
}
```

```
void gen_perm(int n,int m){
    int a[MAXN],temp[MAXN],tag[MAXN]={0},i;
    for (i=0;i<n;i++)
        a[i]=i+1;
    _gen_perm(a,n,m,0,temp,tag);
}
void _gen_comb(int* a,int s,int e,int m,int& count,int* temp){
    int i;
    if (!m)
        dummy(temp,count);
    else
        for (i=s; i \le e-m+1; i++){
            temp[count++]=a[i];
            _gen_comb(a,i+1,e,m-1,count,temp);
            count - -;
        }
}
void gen_comb(int n,int m){
    int a[MAXN],temp[MAXN],count=0,i;
    for (i=0;i<n;i++)
        a[i]=i+1;
   _gen_comb(a,0,n-1,m,count,temp);
}
void _gen_perm_swap(int* a,int n,int l,int* pos,int* dir){
    int i,p1,p2,t;
    if (l==n)
        dummy(a,n);
    else{
        _gen_perm_swap(a,n,l+1,pos,dir);
        for (i=0;i<l;i++){
            p2=(p1=pos[l])+dir[l];
            t=a[p1],a[p1]=a[p2],a[p2]=t;
            pos[a[p1]-1]=p1,pos[a[p2]-1]=p2;
            _gen_perm_swap(a,n,l+1,pos,dir);
        dir[l]=-dir[l];
    }
}
void gen_perm_swap(int n){
    int a[MAXN],pos[MAXN],dir[MAXN],i;
    for (i=0;i<n;i++)
        a[i]=i+1,pos[i]=i,dir[i]=-1;
```

```
_gen_perm_swap(a,n,0,pos,dir);
}
                                  2. 生成 gray 码
//生成 reflected gray code
//每次调用 gray 取得下一个码
//000...000 是第一个码,100...000 是最后一个码
void gray(int n,int *code){
   int t=0,i;
   for (i=0;i<n;t+=code[i++]);</pre>
   if (t&1)
       for (n--;!code[n];n--);
   code[n-1]=1-code[n-1];
}
                                  3. 置换(polya)
//求置换的循环节,polya 原理
//perm[0..n-1]为 0..n-1 的一个置换(排列)
//返回置换最小周期, num 返回循环节个数
#define MAXN 1000
int gcd(int a,int b){
   return b?gcd(b,a%b):a;
}
int polya(int* perm,int n,int& num){
   int i,j,p,v[MAXN]={0},ret=1;
   for (num=i=0;i<n;i++)</pre>
       if (!v[i]){
           for (num++,j=0,p=i;!v[p=perm[p]];j++)
               v[p]=1;
           ret*=j/gcd(ret,j);
       }
   return ret;
}
                                  4. 字典序全排列
//字典序全排列与序号的转换
int perm2num(int n,int *p){
   int i,j,ret=0,k=1;
   for (i=n-2; i>=0; k*=n-(i--))
       for (j=i+1;j<n;j++)
           if (p[j]<p[i])</pre>
               ret+=k;
   return ret;
```

```
}
void num2perm(int n,int *p,int t){
    int i,j;
    for (i=n-1;i>=0;i--)
       p[i]=t%(n-i),t/=n-i;
   for (i=n-1;i;i--)
       for (j=i-1;j>=0;j--)
           if (p[j]<=p[i])</pre>
               p[i]++;
}
                                    5. 字典序组合
//字典序组合与序号的转换
//comb 为组合数 C(n,m),必要时换成大数,注意处理 C(n,m)=0|n<m
int comb(int n,int m){
   int ret=1,i;
   m=m<(n-m)?m:(n-m);
    for (i=n-m+1;i<=n;ret*=(i++));
   for (i=1;i<=m;ret/=(i++));
    return m<0?0:ret;
}
int comb2num(int n,int m,int *c){
    int ret=comb(n,m),i;
    for (i=0;i<m;i++)</pre>
        ret-=comb(n-c[i],m-i);
    return ret;
}
void num2comb(int n,int m,int* c,int t){
    int i,j=1,k;
    for (i=0; i< m; c[i++]=j++)
       for (;t>(k=comb(n-j,m-i-1));t-=k,j++);
}
                                     6. 组合公式
1. C(m,n)=C(m,m-n)
2. C(m,n)=C(m-1,n)+C(m-1,n-1)
derangement D(n) = n!(1 - 1/1! + 1/2! - 1/3! + ... + (-1)^n/n!)
              = (n-1)(D(n-2) - D(n-1))
         Q(n) = D(n) + D(n-1)
求和公式,k = 1...n
1. sum(k) = n(n+1)/2
```

```
2. sum(2k-1) = n^2

3. sum(k^2) = n(n+1)(2n+1)/6

4. sum((2k-1)^2) = n(4n^2-1)/3

5. sum(k^3) = (n(n+1)/2)^2

6. sum((2k-1)^3) = n^2(2n^2-1)

7. sum(k^4) = n(n+1)(2n+1)(3n^2+3n-1)/30

8. sum(k^5) = n^2(n+1)^2(2n^2+2n-1)/12

9. sum(k(k+1)) = n(n+1)(n+2)/3

10. sum(k(k+1)(k+2)) = n(n+1)(n+2)(n+3)/4

12. sum(k(k+1)(k+2)(k+3)) = n(n+1)(n+2)(n+3)(n+4)/5
```

十. 数值计算

1. 定积分计算(Romberg)

```
/* Romberg 求定积分
   输入: 积分区间[a,b], 被积函数 f(x,y,z)
   输出: 积分结果
   f(x,y,z)示例:
   double f0( double x, double l, double t )
       return sqrt(1.0+l*l*t*t*cos(t*x)*cos(t*x));
   }
*/
double Integral (double a, double b, double (*f) (double x, double y, double z), double eps,
               double l, double t)
double Romberg (double a, double b, double (*f) (double x, double y, double z), double eps,
               double l, double t)
{
#define MAX N 1000
    int i, j, temp2, min;
    double h, R[2][MAX_N], temp4;
    for (i=0; i<MAX_N; i++) {
       R[0][i] = 0.0;
       R[1][i] = 0.0;
    }
    h = b-a;
   min = (int)(log(h*10.0)/log(2.0)); //h should be at most 0.1
   R[0][0] = ((*f)(a, l, t)+(*f)(b, l, t))*h*0.50;
    i = 1;
```

```
temp2 = 1;
   while (i<MAX_N){
       i++;
       R[1][0] = 0.0;
       for (j=1; j \le temp2; j++)
           R[1][0] += (*f)(a+h*((double)j-0.50), l, t);
       R[1][0] = (R[0][0] + h*R[1][0])*0.50;
       temp4 = 4.0;
       for (j=1; j<i; j++) {
           R[1][j] = R[1][j-1] + (R[1][j-1]-R[0][j-1])/(temp4-1.0);
           temp4 *= 4.0;
       }
       if ((fabs(R[1][i-1]-R[0][i-2]) < eps) & (i>min))
           return R[1][i-1];
       h *= 0.50;
       temp2 *= 2;
       for (j=0; j<i; j++)
           R[0][j] = R[1][j];
   return R[1][MAX_N-1];
}
double Integral (double a, double b, double (*f) (double x, double y, double z), double eps,
               double l, double t)
#define pi 3.1415926535897932
   int n;
   double R, p, res;
   n = (int)(floor)(b * t * 0.50 / pi);
   p = 2.0 * pi / t;
   res = b - (double)n * p;
   if (n)
       R = Romberg (a, p, f0, eps/(double)n, l, t);
   R = R * (double)n + Romberg( 0.0, res, f0, eps, l, t);
   return R/100.0;
}
                               2. 多项式求根(牛顿法)
/* 牛顿法解多项式的根
  输入: 多项式系数 c[], 多项式度数 n, 求在[a,b]间的根
  输出:根
  要求保证[a,b]间有根
*/
```

```
double fabs( double x )
    return (x<0)? -x : x;
}
double f(int m, double c[], double x)
{
    int i;
    double p = c[m];
    for (i=m; i>0; i--)
        p = p*x + c[i-1];
    return p;
}
int newton(double x0, double *r,
          double c[], double cp[], int n,
          double a, double b, double eps)
{
    int MAX_ITERATION = 1000;
    int i = 1;
    double x1, x2, fp, eps2 = eps/10.0;
   x1 = x0;
   while (i < MAX_ITERATION) {</pre>
       x2 = f(n, c, x1);
        fp = f(n-1, cp, x1);
        if ((fabs(fp)<0.000000001) && (fabs(x2)>1.0))
            return 0;
        x2 = x1 - x2/fp;
        if (fabs(x1-x2) < eps2) {
            if (x2<a || x2>b)
                return 0;
            *r = x2;
            return 1;
        }
        x1 = x2;
        i++;
    }
    return 0;
}
double Polynomial_Root(double c[], int n, double a, double b, double eps)
{
    double *cp;
```

```
int i;
    double root;
    cp = (double *)calloc(n, sizeof(double));
    for (i=n-1; i>=0; i--) {
        cp[i] = (i+1)*c[i+1];
    }
    if (a>b) {
            root = a; a = b; b = root;
    }
    if ((!newton(a, &root, c, cp, n, a, b, eps)) &&
        (!newton(b, &root, c, cp, n, a, b, eps)))
            newton((a+b)*0.5, &root, c, cp, n, a, b, eps);
    free(cp);
    if (fabs(root)<eps)</pre>
        return fabs(root);
    else
        return root;
}
```

3. 周期性方程(追赶法)

```
/* 追赶法解周期性方程
   周期性方程定义: | a1 b1 c1 ...
                                       | | = x1
              | a2 b2 c2 ...
                                        | * | X | = ...
                               an-1 bn-1 | | = xn-1
               | cn-1 ...
               | bn cn
                                an | | = xn
  输入: a[],b[],c[],x[]
  输出: 求解结果 X 在 x[]中
*/
void run()
{
   c[0] /= b[0]; a[0] /= b[0]; x[0] /= b[0];
   for (int i = 1; i < N - 1; i ++) {
      double temp = b[i] - a[i] * c[i - 1];
       c[i] /= temp;
      x[i] = (x[i] - a[i] * x[i - 1]) / temp;
      a[i] = -a[i] * a[i - 1] / temp;
   }
   a[N - 2] = -a[N - 2] - c[N - 2];
   for (int i = N - 3; i >= 0; i --) {
       a[i] = -a[i] - c[i] * a[i + 1];
```

十一. 几何

1. 多边形

```
#include <stdlib.h>
#include <math.h>
#define MAXN 1000
#define offset 10000
#define eps 1e-8
#define zero(x) (((x)>0?(x):-(x))<eps)
#define _sign(x) ((x)>eps?1:((x)<-eps?2:0))
struct point{double x,y;};
struct line{point a,b;};
double xmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
//判定凸多边形,顶点按顺时针或逆时针给出,允许相邻边共线
int is_convex(int n,point* p){
   int i,s[3]=\{1,1,1\};
   for (i=0; i<n\&\&s[1]|s[2]; i++)
       s[_sign(xmult(p[(i+1)%n],p[(i+2)%n],p[i]))]=0;
   return s[1]|s[2];
}
//判定凸多边形,顶点按顺时针或逆时针给出,不允许相邻边共线
int is_convex_v2(int n,point* p){
   int i,s[3]=\{1,1,1\};
   for (i=0; i<n\&\&s[0]\&\&s[1]|s[2]; i++)
       s[_sign(xmult(p[(i+1)%n],p[(i+2)%n],p[i]))]=0;
   return s[0]&&s[1]|s[2];
}
//判点在凸多边形内或多边形边上,顶点按顺时针或逆时针给出
int inside_convex(point q,int n,point* p){
   int i,s[3]=\{1,1,1\};
```

```
for (i=0; i<n\&\&s[1]|s[2]; i++)
                    s[_sign(xmult(p[(i+1)%n],q,p[i]))]=0;
          return s[1]|s[2];
}
//判点在凸多边形内,顶点按顺时针或逆时针给出,在多边形边上返回 0
int inside_convex_v2(point q,int n,point* p){
          int i,s[3]=\{1,1,1\};
          for (i=0; i<n\&\&s[0]\&\&s[1]|s[2]; i++)
                   s[_sign(xmult(p[(i+1)%n],q,p[i]))]=0;
          return s[0]&&s[1]|s[2];
}
//判点在任意多边形内,顶点按顺时针或逆时针给出
//on_edge 表示点在多边形边上时的返回值,offset 为多边形坐标上限
int inside_polygon(point q,int n,point* p,int on_edge=1){
          point q2;
          int i=0,count;
         while (i<n)
                    for (count=i=0,q2.x=rand()+offset,q2.y=rand()+offset;i<n;i++)</pre>
                              if
(zero(xmult(q,p[i],p[(i+1)%n]))\&\&(p[i].x-q.x)*(p[(i+1)%n].x-q.x)<eps\&\&(p[i].y-q.y)*(p[(i+1)%n])
+1)%n].y-q.y)<eps)
                                       return on_edge;
                             else if (zero(xmult(q,q2,p[i])))
                                       break;
                             else
                                                                                                                                                                                                                          if
(xmult(q,p[i],q2)*xmult(q,p[(i+1)%n],q2)<-eps&&xmult(p[i],q,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%n])*xmult(p[i],q2,p[(i+1)%
(i+1)%n])<-eps)
                                       count++;
          return count&1;
}
inline int opposite side(point p1, point p2, point l1, point l2){
          return xmult(l1,p1,l2)*xmult(l1,p2,l2)<-eps;</pre>
}
inline int dot_online_in(point p,point l1,point l2){
          return zero(xmult(p,l1,l2))&&(l1.x-p.x)*(l2.x-p.x)<eps&&(l1.y-p.y)*(l2.y-p.y)<eps;
}
//判线段在任意多边形内,顶点按顺时针或逆时针给出,与边界相交返回1
int inside polygon(point l1,point l2,int n,point* p){
          point t[MAXN],tt;
          int i, j, k=0;
          if (!inside_polygon(l1,n,p)||!inside_polygon(l2,n,p))
```

```
return 0;
    for (i=0;i<n;i++)
        if (opposite\_side(l1,l2,p[i],p[(i+1)%n])\&&opposite\_side(p[i],p[(i+1)%n],l1,l2))
        else if (dot_online_in(l1,p[i],p[(i+1)%n]))
            t[k++]=l1;
        else if (dot_online_in(l2,p[i],p[(i+1)%n]))
            t[k++]=l2;
        else if (dot_online_in(p[i],l1,l2))
            t[k++]=p[i];
    for (i=0;i<k;i++)
        for (j=i+1; j< k; j++){
            tt.x=(t[i].x+t[j].x)/2;
           tt.y=(t[i].y+t[j].y)/2;
            if (!inside_polygon(tt,n,p))
                return 0;
        }
    return 1;
}
point intersection(line u,line v){
    point ret=u.a;
    double t=((u.a.x-v.a.x)*(v.a.y-v.b.y)-(u.a.y-v.a.y)*(v.a.x-v.b.x))
            /((u.a.x-u.b.x)*(v.a.y-v.b.y)-(u.a.y-u.b.y)*(v.a.x-v.b.x));
    ret.x+=(u.b.x-u.a.x)*t;
    ret.y+=(u.b.y-u.a.y)*t;
    return ret;
}
point barycenter(point a,point b,point c){
   line u,v;
    u.a.x=(a.x+b.x)/2;
   u.a.y=(a.y+b.y)/2;
   u.b=c;
    v.a.x=(a.x+c.x)/2;
   v.a.y=(a.y+c.y)/2;
   v.b=b;
    return intersection(u,v);
}
//多边形重心
point barycenter(int n,point* p){
    point ret,t;
    double t1=0,t2;
    int i;
    ret.x=ret.y=0;
```

```
for (i=1;i<n-1;i++)
       if (fabs(t2=xmult(p[0],p[i],p[i+1]))>eps){
           t=barycenter(p[0],p[i],p[i+1]);
            ret.x+=t.x*t2;
           ret.y+=t.y*t2;
           t1+=t2:
       }
    if (fabs(t1)>eps)
        ret.x/=t1, ret.y/=t1;
    return ret;
}
                                     2. 多边形切割
//多边形切割
//可用于半平面交
#define MAXN 100
#define eps 1e-8
#define zero(x) (((x)>0?(x):-(x))<eps)
struct point{double x,y;};
double xmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
int same_side(point p1,point p2,point l1,point l2){
    return xmult(l1,p1,l2)*xmult(l1,p2,l2)>eps;
}
point intersection(point u1,point u2,point v1,point v2){
    point ret=u1;
    double t=((u1.x-v1.x)*(v1.y-v2.y)-(u1.y-v1.y)*(v1.x-v2.x))
           /((u1.x-u2.x)*(v1.y-v2.y)-(u1.y-u2.y)*(v1.x-v2.x));
    ret.x+=(u2.x-u1.x)*t;
    ret.y+=(u2.y-u1.y)*t;
    return ret;
}
//将多边形沿 l1, l2 确定的直线切割在 side 侧切割, 保证 l1, l2, side 不共线
void polygon_cut(int& n,point* p,point l1,point l2,point side){
    point pp[100];
    int m=0,i;
    for (i=0; i< n; i++){
       if (same_side(p[i],side,l1,l2))
            pp[m++]=p[i];
       if
(!same\_side(p[i],p[(i+1)%n],l1,l2)&\&!(zero(xmult(p[i],l1,l2))&&zero(xmult(p[(i+1)%n],l1,l2))&
```

```
12))))
           pp[m++]=intersection(p[i],p[(i+1)%n],l1,l2);
    }
    for (n=i=0;i< m;i++)
       if (!i||!zero(pp[i].x-pp[i-1].x)||!zero(pp[i].y-pp[i-1].y))
           p[n++]=pp[i];
    if (zero(p[n-1].x-p[0].x)&&zero(p[n-1].y-p[0].y))
       n - - :
    if (n<3)
       n=0;
}
                                      3. 浮点函数
//浮点几何函数库
#include <math.h>
#define eps 1e-8
#define zero(x) (((x)>0?(x):-(x))<eps)
struct point{double x,y;};
struct line{point a,b;};
//计算 cross product (P1-P0)x(P2-P0)
double xmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
double xmult(double x1,double y1,double x2,double y2,double x0,double y0){
    return (x1-x0)*(y2-y0)-(x2-x0)*(y1-y0);
}
//计算 dot product (P1-P0).(P2-P0)
double dmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.x-p0.x)+(p1.y-p0.y)*(p2.y-p0.y);
}
double dmult(double x1,double y1,double x2,double y2,double x0,double y0){
    return (x1-x0)*(x2-x0)+(y1-y0)*(y2-y0);
}
//两点距离
double distance(point p1,point p2){
    return sqrt((p1.x-p2.x)*(p1.x-p2.x)+(p1.y-p2.y)*(p1.y-p2.y));
}
double distance(double x1,double y1,double x2,double y2){
    return sqrt((x1-x2)*(x1-x2)+(y1-y2)*(y1-y2));
}
//判三点共线
int dots_inline(point p1,point p2,point p3){
```

```
return zero(xmult(p1,p2,p3));
}
int dots inline(double x1,double y1,double x2,double y2,double x3,double y3){
   return zero(xmult(x1,y1,x2,y2,x3,y3));
}
//判点是否在线段上,包括端点
int dot_online_in(point p,line l){
zero(xmult(p,l.a,l.b))&&(l.a.x-p.x)*(l.b.x-p.x)<eps&&(l.a.y-p.y)*(l.b.y-p.y)<eps;
}
int dot_online_in(point p,point l1,point l2){
   return \ zero(xmult(p,l1,l2))\&\&(l1.x-p.x)*(l2.x-p.x) < eps\&&(l1.y-p.y)*(l2.y-p.y) < eps;\\
}
int dot_online_in(double x,double y,double x1,double y1,double x2,double y2){
   return zero(xmult(x,y,x1,y1,x2,y2))&(x1-x)*(x2-x)<eps&(y1-y)*(y2-y)<eps;
}
//判点是否在线段上,不包括端点
int dot_online_ex(point p,line l){
   return
dot\_online\_in(p,l)\&\&(!zero(p.x-l.a.x)||!zero(p.y-l.a.y))\&\&(!zero(p.x-l.b.x)||!zero(p.y-l.a.y)|
}
int dot_online_ex(point p,point l1,point l2){
l2.y));
int dot_online_ex(double x,double y,double x1,double y1,double x2,double y2){
   return
dot_online_in(x,y,x1,y1,x2,y2)\&(!zero(x-x1)||!zero(y-y1))\&(!zero(x-x2)||!zero(y-y2));
//判两点在线段同侧,点在线段上返回 0
int same side(point p1,point p2,line l){
   return xmult(l.a,p1,l.b)*xmult(l.a,p2,l.b)>eps;
}
int same_side(point p1,point p2,point l1,point l2){
   return xmult(l1,p1,l2)*xmult(l1,p2,l2)>eps;
}
//判两点在线段异侧,点在线段上返回 0
int opposite_side(point p1,point p2,line l){
   return xmult(l.a,p1,l.b)*xmult(l.a,p2,l.b)<-eps;</pre>
}
```

```
int opposite_side(point p1,point p2,point l1,point l2){
    return xmult(l1,p1,l2)*xmult(l1,p2,l2)<-eps;</pre>
}
// 点关于直线的对称点 // by lyt
// 缺点: 用了斜率
// 也可以利用"点到直线上的最近点"来做,避免使用斜率。
point symmetric point(point p1, point l1, point l2) {
 if (l1.x > l2.x - eps \&\& l1.x < l2.x + eps) {
   ret.x = (2 * l1.x - p1.x);
   ret.y = pl.y;
 } else {
   double k = (l1.y - l2.y) / (l1.x - l2.x);
   ret.x = (2*k*k*l1.x + 2*k*p1.y - 2*k*l1.y - k*k*p1.x + p1.x) / (1 + k*k);
   ret.y = p1.y - (ret.x - p1.x) / k;
 }
 return ret;
}
//判两直线平行
int parallel(line u,line v){
    return zero((u.a.x-u.b.x)*(v.a.y-v.b.y)-(v.a.x-v.b.x)*(u.a.y-u.b.y));
}
int parallel(point u1,point u2,point v1,point v2){
    return zero((u1.x-u2.x)*(v1.y-v2.y)-(v1.x-v2.x)*(u1.y-u2.y));
}
//判两直线垂直
int perpendicular(line u,line v){
    return zero((u.a.x-u.b.x)*(v.a.x-v.b.x)+(u.a.y-u.b.y)*(v.a.y-v.b.y));
}
int perpendicular(point u1,point u2,point v1,point v2){
    return zero((u1.x-u2.x)*(v1.x-v2.x)+(u1.y-u2.y)*(v1.y-v2.y));
}
//判两线段相交,包括端点和部分重合
int intersect_in(line u,line v){
    if (!dots_inline(u.a,u.b,v.a)||!dots_inline(u.a,u.b,v.b))
       return \ !same\_side(u.a,u.b,v)\&\&!same\_side(v.a,v.b,u);\\
    return
dot_online_in(u.a,v)||dot_online_in(u.b,v)||dot_online_in(v.a,u)||dot_online_in(v.b,u);
int intersect_in(point u1,point u2,point v1,point v2){
    if (!dots_inline(u1,u2,v1)||!dots_inline(u1,u2,v2))
        return !same_side(u1,u2,v1,v2)&&!same_side(v1,v2,u1,u2);
```

```
return
dot_online_in(u1,v1,v2)||dot_online_in(u2,v1,v2)||dot_online_in(v1,u1,u2)||dot_online_in
(v2,u1,u2);
}
//判两线段相交,不包括端点和部分重合
int intersect_ex(line u,line v){
   return opposite_side(u.a,u.b,v)&&opposite_side(v.a,v.b,u);
}
int intersect_ex(point u1,point u2,point v1,point v2){
   return opposite_side(u1,u2,v1,v2)&&opposite_side(v1,v2,u1,u2);
}
//计算两直线交点,注意事先判断直线是否平行!
//线段交点请另外判线段相交(同时还是要判断是否平行!)
point intersection(line u,line v){
   point ret=u.a;
   double t=((u.a.x-v.a.x)*(v.a.y-v.b.y)-(u.a.y-v.a.y)*(v.a.x-v.b.x))
           /((u.a.x-u.b.x)*(v.a.y-v.b.y)-(u.a.y-u.b.y)*(v.a.x-v.b.x));
   ret.x+=(u.b.x-u.a.x)*t;
   ret.y+=(u.b.y-u.a.y)*t;
   return ret:
}
point intersection(point u1,point u2,point v1,point v2){
   point ret=u1;
   double t=((u1.x-v1.x)*(v1.y-v2.y)-(u1.y-v1.y)*(v1.x-v2.x))
           /((u1.x-u2.x)*(v1.y-v2.y)-(u1.y-u2.y)*(v1.x-v2.x));
   ret.x+=(u2.x-u1.x)*t;
   ret.y+=(u2.y-u1.y)*t;
   return ret;
}
//点到直线上的最近点
point ptoline(point p,line l){
   point t=p;
   t.x+=l.a.y-l.b.y,t.y+=l.b.x-l.a.x;
   return intersection(p,t,l.a,l.b);
point ptoline(point p,point l1,point l2){
   point t=p;
   t.x+=l1.y-l2.y,t.y+=l2.x-l1.x;
   return intersection(p,t,l1,l2);
}
//点到直线距离
```

double disptoline(point p,line l){

```
return fabs(xmult(p,l.a,l.b))/distance(l.a,l.b);
}
double disptoline(point p,point l1,point l2){
    return fabs(xmult(p,l1,l2))/distance(l1,l2);
}
double disptoline(double x,double y,double x1,double y1,double x2,double y2){
    return fabs(xmult(x,y,x1,y1,x2,y2))/distance(x1,y1,x2,y2);
}
//点到线段上的最近点
point ptoseg(point p,line l){
   point t=p;
   t.x+=l.a.y-l.b.y,t.y+=l.b.x-l.a.x;
   if (xmult(l.a,t,p)*xmult(l.b,t,p)>eps)
        return distance(p,l.a)<distance(p,l.b)?l.a:l.b;</pre>
    return intersection(p,t,l.a,l.b);
}
point ptoseg(point p,point l1,point l2){
   point t=p;
   t.x+=l1.y-l2.y, t.y+=l2.x-l1.x;
   if (xmult(l1,t,p)*xmult(l2,t,p)>eps)
        return distance(p,l1)<distance(p,l2)?l1:l2;</pre>
   return intersection(p,t,l1,l2);
}
//点到线段距离
double disptoseg(point p,line l){
   point t=p;
   t.x+=l.a.y-l.b.y,t.y+=l.b.x-l.a.x;
   if (xmult(l.a,t,p)*xmult(l.b,t,p)>eps)
        return distance(p,l.a)<distance(p,l.b)?distance(p,l.a):distance(p,l.b);</pre>
    return fabs(xmult(p,l.a,l.b))/distance(l.a,l.b);
}
double disptoseg(point p,point l1,point l2){
   point t=p;
   t.x+=l1.y-l2.y, t.y+=l2.x-l1.x;
   if (xmult(l1,t,p)*xmult(l2,t,p)>eps)
        return distance(p,l1)<distance(p,l2)?distance(p,l1):distance(p,l2);</pre>
    return fabs(xmult(p,l1,l2))/distance(l1,l2);
}
//矢量 V 以 P 为顶点逆时针旋转 angle 并放大 scale 倍
point rotate(point v,point p,double angle,double scale){
   point ret=p;
   v.x-=p.x, v.y-=p.y;
   p.x=scale*cos(angle);
```

```
p.y=scale*sin(angle);
ret.x+=v.x*p.x-v.y*p.y;
ret.y+=v.x*p.y+v.y*p.x;
return ret;
}
```

4. 几何公式

三角形:

- 1. 半周长 P=(a+b+c)/2
- 2. 面积 S=aHa/2=absin(C)/2=sqrt(P(P-a)(P-b)(P-c))
- 3. 中线 Ma=sqrt(2(b^2+c^2)-a^2)/2=sqrt(b^2+c^2+2bccos(A))/2
- 4. 角平分线 Ta=sqrt(bc((b+c)^2-a^2))/(b+c)=2bccos(A/2)/(b+c)
- 5. 高线 Ha=bsin(C)=csin(B)=sqrt(b^2-((a^2+b^2-c^2)/(2a))^2)
- 6. 内切圆半径 r=S/P=asin(B/2)sin(C/2)/sin((B+C)/2)=4Rsin(A/2)sin(B/2)sin(C/2)=sqrt((P-a)(P-b)(P-c)/P)

=Ptan(A/2)tan(B/2)tan(C/2)

7. 外接圆半径 R=abc/(4S)=a/(2sin(A))=b/(2sin(B))=c/(2sin(C))

四边形:

D1,D2 为对角线,M对角线中点连线,A为对角线夹角

- 1. a^2+b^2+c^2+d^2=D1^2+D2^2+4M^2
- 2. S=D1D2sin(A)/2

(以下对圆的内接四边形)

- 3. ac+bd=D1D2
- 4. S=sqrt((P-a)(P-b)(P-c)(P-d)),P 为半周长

正 n 边形:

R 为外接圆半径, r 为内切圆半径

- 1. 中心角 A=2PI/n
- 2. 内角 C=(n-2)PI/n
- 3. 边长 a=2sqrt(R^2-r^2)=2Rsin(A/2)=2rtan(A/2)
- 4. 面积 S=nar/2=nr^2tan(A/2)=nR^2sin(A)/2=na^2/(4tan(A/2))

圆:

- 1. 弧长 l=rA
- 2. 弦长 a=2sqrt(2hr-h^2)=2rsin(A/2)
- 3. 弓形高 h=r-sqrt(r^2-a^2/4)=r(1-cos(A/2))=atan(A/4)/2
- 4. 扇形面积 S1=rl/2=r^2A/2
- 5. 弓形面积 S2=(rl-a(r-h))/2=r^2(A-sin(A))/2

棱柱:

- 1. 体积 V=Ah,A 为底面积,h 为高
- 2. 侧面积 S=lp,l 为棱长,p 为直截面周长
- 3. 全面积 T=S+2A

棱锥:

- 1. 体积 V=Ah/3,A 为底面积,h 为高
- (以下对正棱锥)
- 2. 侧面积 S=lp/2,l 为斜高,p 为底面周长
- 3. 全面积 T=S+A

棱台:

- 1. 体积 V=(A1+A2+sqrt(A1A2))h/3,A1.A2 为上下底面积,h 为高(以下为正棱台)
- 2. 侧面积 S=(p1+p2)1/2,p1.p2 为上下底面周长,l 为斜高
- 3. 全面积 T=S+A1+A2

圆柱:

- 1. 侧面积 S=2PIrh
- 2. 全面积 T=2PIr(h+r)
- 3. 体积 V=PIr^2h

圆锥:

- 1. 母线 l=sqrt(h^2+r^2)
- 2. 侧面积 S=PIrl
- 3. 全面积 T=PIr(l+r)
- 4. 体积 V=PIr^2h/3

圆台:

- 1. 母线 l=sqrt(h^2+(r1-r2)^2)
- 2. 侧面积 S=PI(r1+r2)l
- 3. 全面积 T=PIr1(l+r1)+PIr2(l+r2)
- 4. 体积 V=PI(r1^2+r2^2+r1r2)h/3

球:

- 1. 全面积 T=4PIr^2
- 2. 体积 V=4PIr^3/3

球台:

- 1. 侧面积 S=2PIrh
- 2. 全面积 T=PI(2rh+r1^2+r2^2)
- 3. 体积 V=PIh(3(r1^2+r2^2)+h^2)/6

球扇形:

- 1. 全面积 T=PIr(2h+r0), h 为球冠高, r0 为球冠底面半径
- 2. 体积 V=2PIr^2h/3

5. 面积

#include <math.h>
struct point{double x,y;};

//计算 cross product (P1-P0)x(P2-P0)

```
double xmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
double xmult(double x1,double y1,double x2,double y2,double x0,double y0){
    return (x1-x0)*(y2-y0)-(x2-x0)*(y1-y0);
}
//计算三角形面积,输入三顶点
double area_triangle(point p1,point p2,point p3){
    return fabs(xmult(p1,p2,p3))/2;
}
double area_triangle(double x1,double y1,double x2,double y2,double x3,double y3){
    return fabs(xmult(x1,y1,x2,y2,x3,y3))/2;
}
//计算三角形面积,输入三边长
double area_triangle(double a,double b,double c){
   double s=(a+b+c)/2;
   return sqrt(s*(s-a)*(s-b)*(s-c));
}
//计算多边形面积,顶点按顺时针或逆时针给出
double area_polygon(int n,point* p){
   double s1=0, s2=0;
   int i;
   for (i=0;i<n;i++)
       s1+=p[(i+1)%n].y*p[i].x,s2+=p[(i+1)%n].y*p[(i+2)%n].x;
   return fabs(s1-s2)/2;
}
                                       6. 球面
#include <math.h>
const double pi=acos(-1);
//计算圆心角 lat 表示纬度, -90<=w<=90, lng 表示经度
//返回两点所在大圆劣弧对应圆心角,0<=angle<=pi
double angle(double lng1,double lat1,double lng2,double lat2){
   double dlng=fabs(lng1-lng2)*pi/180;
   while (dlng>=pi+pi)
       dlng-=pi+pi;
   if (dlng>pi)
       dlng=pi+pi-dlng;
   lat1*=pi/180,lat2*=pi/180;
    return acos(cos(lat1)*cos(lat2)*cos(dlng)+sin(lat1)*sin(lat2));
}
```

```
//计算距离,r 为球半径
double line_dist(double r,double lng1,double lat1,double lng2,double lat2){
    double dlng=fabs(lng1-lng2)*pi/180;
   while (dlng>=pi+pi)
       dlng-=pi+pi;
    if (dlng>pi)
       dlng=pi+pi-dlng;
   lat1*=pi/180,lat2*=pi/180;
    return r*sqrt(2-2*(cos(lat1)*cos(lat2)*cos(dlng)+sin(lat1)*sin(lat2)));
}
//计算球面距离,r为球半径
inline double sphere dist(double r,double lng1,double lat1,double lng2,double lat2){
    return r*angle(lng1,lat1,lng2,lat2);
}
                                       7. 三角形
#include <math.h>
struct point{double x,y;};
struct line{point a,b;};
double distance(point p1,point p2){
    return sqrt((p1.x-p2.x)*(p1.x-p2.x)+(p1.y-p2.y)*(p1.y-p2.y));
}
point intersection(line u,line v){
    point ret=u.a;
    double t=((u.a.x-v.a.x)*(v.a.y-v.b.y)-(u.a.y-v.a.y)*(v.a.x-v.b.x))
           /((u.a.x-u.b.x)*(v.a.y-v.b.y)-(u.a.y-u.b.y)*(v.a.x-v.b.x));
    ret.x+=(u.b.x-u.a.x)*t;
    ret.y+=(u.b.y-u.a.y)*t;
    return ret;
}
//外心
point circumcenter(point a,point b,point c){
    line u,v;
    u.a.x=(a.x+b.x)/2;
    u.a.y=(a.y+b.y)/2;
    u.b.x=u.a.x-a.y+b.y;
    u.b.y=u.a.y+a.x-b.x;
    v.a.x=(a.x+c.x)/2;
    v.a.y=(a.y+c.y)/2;
    v.b.x=v.a.x-a.y+c.y;
    v.b.y=v.a.y+a.x-c.x;
    return intersection(u,v);
```

```
}
//内心
point incenter(point a,point b,point c){
   line u,v;
   double m,n;
   u.a=a;
   m=atan2(b.y-a.y,b.x-a.x);
   n=atan2(c.y-a.y,c.x-a.x);
   u.b.x=u.a.x+cos((m+n)/2);
   u.b.y=u.a.y+sin((m+n)/2);
   v.a=b;
   m=atan2(a.y-b.y,a.x-b.x);
   n=atan2(c.y-b.y,c.x-b.x);
   v.b.x=v.a.x+cos((m+n)/2);
   v.b.y=v.a.y+sin((m+n)/2);
   return intersection(u,v);
}
//垂心
point perpencenter(point a,point b,point c){
   line u,v;
   u.a=c;
   u.b.x=u.a.x-a.y+b.y;
   u.b.y=u.a.y+a.x-b.x;
   v.a=b;
   v.b.x=v.a.x-a.y+c.y;
   v.b.y=v.a.y+a.x-c.x;
   return intersection(u,v);
}
//重心
//到三角形三顶点距离的平方和最小的点
//三角形内到三边距离之积最大的点
point barycenter(point a,point b,point c){
   line u,v;
   u.a.x=(a.x+b.x)/2;
   u.a.y=(a.y+b.y)/2;
   u.b=c;
   v.a.x=(a.x+c.x)/2;
   v.a.y=(a.y+c.y)/2;
   v.b=b;
   return intersection(u,v);
}
```

```
//到三角形三顶点距离之和最小的点
point fermentpoint(point a,point b,point c){
    point u,v;
    double step=fabs(a.x)+fabs(a.y)+fabs(b.x)+fabs(b.y)+fabs(c.x)+fabs(c.y);
    int i,j,k;
    u.x=(a.x+b.x+c.x)/3;
    u.y=(a.y+b.y+c.y)/3;
   while (step>1e-10)
       for (k=0; k<10; step/=2, k++)
           for (i=-1;i<=1;i++)
               for (j=-1; j<=1; j++){
                   v.x=u.x+step*i;
                   v.y=u.y+step*j;
                   if
(distance(u,a)+distance(u,b)+distance(u,c)>distance(v,a)+distance(v,b)+distance(v,c))
                       u=v;
               }
    return u;
}
                                     8. 三维几何
//三维几何函数库
#include <math.h>
#define eps 1e-8
#define zero(x) (((x)>0?(x):-(x))<eps)
struct point3{double x,y,z;};
struct line3{point3 a,b;};
struct plane3{point3 a,b,c;};
//计算 cross product U x V
point3 xmult(point3 u,point3 v){
    point3 ret;
    ret.x=u.y*v.z-v.y*u.z;
    ret.y=u.z*v.x-u.x*v.z;
    ret.z=u.x*v.y-u.y*v.x;
    return ret;
}
//计算 dot product U . V
double dmult(point3 u,point3 v){
    return u.x*v.x+u.y*v.y+u.z*v.z;
}
//矢量差 U - V
point3 subt(point3 u,point3 v){
    point3 ret;
```

```
ret.x=u.x-v.x;
     ret.y=u.y-v.y;
     ret.z=u.z-v.z;
     return ret;
}
//取平面法向量
point3 pvec(plane3 s){
     return xmult(subt(s.a,s.b),subt(s.b,s.c));
}
point3 pvec(point3 s1,point3 s2,point3 s3){
     return xmult(subt(s1,s2),subt(s2,s3));
}
//两点距离,单参数取向量大小
double distance(point3 p1,point3 p2){
     return
sqrt((p1.x-p2.x)*(p1.x-p2.x)+(p1.y-p2.y)*(p1.y-p2.y)+(p1.z-p2.z)*(p1.z-p2.z));
}
//向量大小
double vlen(point3 p){
     return sqrt(p.x*p.x+p.y*p.y+p.z*p.z);
}
//判三点共线
int dots inline(point3 p1,point3 p2,point3 p3){
     return vlen(xmult(subt(p1,p2),subt(p2,p3)))<eps;</pre>
}
//判四点共面
int dots_onplane(point3 a,point3 b,point3 c,point3 d){
     return zero(dmult(pvec(a,b,c),subt(d,a)));
}
//判点是否在线段上,包括端点和共线
int dot_online_in(point3 p,line3 l){
     return zero(vlen(xmult(subt(p,l.a),subt(p,l.b))))\&\&(l.a.x-p.x)*(l.b.x-p.x)<eps&&
           (l.a.y-p.y)*(l.b.y-p.y) < eps&(l.a.z-p.z)*(l.b.z-p.z) < eps;
}
int dot_online_in(point3 p,point3 l1,point3 l2){
     return\ zero(vlen(xmult(subt(p,l1),subt(p,l2))))\&\&(l1.x-p.x)*(l2.x-p.x)<eps\&\&line(xmult(subt(p,l1),subt(p,l2))))\&\&(line(xmult(subt(p,l1),subt(p,l2))))\&(line(xmult(subt(p,l1),subt(p,l2))))\&(line(xmult(subt(p,l1),subt(p,l2))))\&(line(xmult(subt(p,l1),subt(p,l2)))))\&(line(xmult(subt(p,l1),subt(p,l2)))))\&(line(xmult(subt(p,l1),subt(p,l2)))))\&(line(xmult(subt(p,l1),subt(p,l2)))))\&(line(xmult(subt(p,l1),subt(p,l2)))))\&(line(xmult(subt(p,l2),subt(p,l2))))))\\
           (l1.y-p.y)*(l2.y-p.y) < eps&&(l1.z-p.z)*(l2.z-p.z) < eps;
}
```

```
int dot_online_ex(point3 p,line3 l){
   return dot_online_in(p,l)\&\&(!zero(p.x-l.a.x)||!zero(p.y-l.a.y)||!zero(p.z-l.a.z))\&\&
       (!zero(p.x-l.b.x)||!zero(p.y-l.b.y)||!zero(p.z-l.b.z));
}
int dot_online_ex(point3 p,point3 l1,point3 l2){
   return dot_online_in(p,l1,l2)\&\&(!zero(p.x-l1.x))||!zero(p.y-l1.y)||!zero(p.z-l1.z))\&\&
       (!zero(p.x-l2.x)||!zero(p.y-l2.y)||!zero(p.z-l2.z));
}
//判点是否在空间三角形上,包括边界,三点共线无意义
int dot_inplane_in(point3 p,plane3 s){
   return
zero(vlen(xmult(subt(s.a,s.b),subt(s.a,s.c))) - vlen(xmult(subt(p,s.a),subt(p,s.b))) -
       vlen(xmult(subt(p,s.b),subt(p,s.c)))-vlen(xmult(subt(p,s.c),subt(p,s.a))));
}
int dot_inplane_in(point3 p,point3 s1,point3 s2,point3 s3){
   return zero(vlen(xmult(subt(s1,s2),subt(s1,s3)))-vlen(xmult(subt(p,s1),subt(p,s2)))-
       vlen(xmult(subt(p,s2),subt(p,s3))) - vlen(xmult(subt(p,s3),subt(p,s1))));
}
//判点是否在空间三角形上,不包括边界,三点共线无意义
int dot_inplane_ex(point3 p,plane3 s){
   return dot_inplane_in(p,s)&&vlen(xmult(subt(p,s.a),subt(p,s.b)))>eps&&
   vlen(xmult(subt(p,s.b),subt(p,s.c))) > eps&vlen(xmult(subt(p,s.c),subt(p,s.a))) > eps;
}
int dot_inplane_ex(point3 p,point3 s1,point3 s2,point3 s3){
   return dot_inplane_in(p,s1,s2,s3)&&vlen(xmult(subt(p,s1),subt(p,s2)))>eps&&
       vlen(xmult(subt(p,s2),subt(p,s3))) > eps&vlen(xmult(subt(p,s3),subt(p,s1))) > eps;
}
//判两点在线段同侧,点在线段上返回 0,不共面无意义
int same_side(point3 p1,point3 p2,line3 l){
   return
dmult(xmult(subt(l.a,l.b)),subt(p1,l.b)),xmult(subt(l.a,l.b)),subt(p2,l.b)))>eps;
int same_side(point3 p1,point3 p2,point3 l1,point3 l2){
   return dmult(xmult(subt(l1,l2),subt(p1,l2)),xmult(subt(l1,l2),subt(p2,l2)))>eps;
}
//判两点在线段异侧,点在线段上返回0,不共面无意义
int opposite_side(point3 p1,point3 p2,line3 l){
    return
dmult(xmult(subt(l.a,l.b),subt(p1,l.b)),xmult(subt(l.a,l.b),subt(p2,l.b))) <-eps;
int opposite_side(point3 p1,point3 p2,point3 l1,point3 l2){
```

```
return dmult(xmult(subt(l1,l2),subt(p1,l2)),xmult(subt(l1,l2),subt(p2,l2)))<-eps;</pre>
}
//判两点在平面同侧,点在平面上返回 0
int same_side(point3 p1,point3 p2,plane3 s){
    return dmult(pvec(s),subt(p1,s.a))*dmult(pvec(s),subt(p2,s.a))>eps;
}
int same_side(point3 p1,point3 p2,point3 s1,point3 s2,point3 s3){
    return dmult(pvec(s1,s2,s3),subt(p1,s1))*dmult(pvec(s1,s2,s3),subt(p2,s1))>eps;
}
//判两点在平面异侧,点在平面上返回 \theta
int opposite_side(point3 p1,point3 p2,plane3 s){
    return dmult(pvec(s),subt(p1,s.a))*dmult(pvec(s),subt(p2,s.a))<-eps;</pre>
}
int opposite_side(point3 p1,point3 p2,point3 s1,point3 s2,point3 s3){
    return dmult(pvec(s1,s2,s3),subt(p1,s1))*dmult(pvec(s1,s2,s3),subt(p2,s1))<-eps;
}
//判两直线平行
int parallel(line3 u,line3 v){
    return vlen(xmult(subt(u.a,u.b),subt(v.a,v.b)))<eps;</pre>
}
int parallel(point3 u1,point3 u2,point3 v1,point3 v2){
    return vlen(xmult(subt(u1,u2),subt(v1,v2)))<eps;</pre>
}
//判两平面平行
int parallel(plane3 u,plane3 v){
    return vlen(xmult(pvec(u),pvec(v)))<eps;</pre>
}
int parallel(point3 u1,point3 u2,point3 u3,point3 v1,point3 v2,point3 v3){
    return vlen(xmult(pvec(u1,u2,u3),pvec(v1,v2,v3)))<eps;</pre>
}
//判直线与平面平行
int parallel(line3 l,plane3 s){
    return zero(dmult(subt(l.a,l.b),pvec(s)));
}
int parallel(point3 l1,point3 l2,point3 s1,point3 s2,point3 s3){
    return zero(dmult(subt(l1,l2),pvec(s1,s2,s3)));
}
//判两直线垂直
int perpendicular(line3 u,line3 v){
    return zero(dmult(subt(u.a,u.b),subt(v.a,v.b)));
```

```
}
int perpendicular(point3 u1,point3 u2,point3 v1,point3 v2){
    return zero(dmult(subt(u1,u2),subt(v1,v2)));
}
//判两平面垂直
int perpendicular(plane3 u,plane3 v){
    return zero(dmult(pvec(u),pvec(v)));
}
int perpendicular(point3 u1,point3 u2,point3 u3,point3 v1,point3 v2,point3 v3){
    return zero(dmult(pvec(u1,u2,u3),pvec(v1,v2,v3)));
}
//判直线与平面平行
int perpendicular(line3 l,plane3 s){
    return vlen(xmult(subt(l.a,l.b),pvec(s)))<eps;</pre>
}
int perpendicular(point3 l1,point3 l2,point3 s1,point3 s2,point3 s3){
    return vlen(xmult(subt(l1,l2),pvec(s1,s2,s3)))<eps;</pre>
}
//判两线段相交,包括端点和部分重合
int intersect_in(line3 u,line3 v){
   if (!dots_onplane(u.a,u.b,v.a,v.b))
        return 0;
   if (!dots_inline(u.a,u.b,v.a)||!dots_inline(u.a,u.b,v.b))
        return !same_side(u.a,u.b,v)&&!same_side(v.a,v.b,u);
    return
dot_online_in(u.a,v)||dot_online_in(u.b,v)||dot_online_in(v.a,u)||dot_online_in(v.b,u);
int intersect_in(point3 u1,point3 u2,point3 v1,point3 v2){
   if (!dots_onplane(u1,u2,v1,v2))
        return 0;
    if (!dots inline(u1,u2,v1)||!dots inline(u1,u2,v2))
       return !same_side(u1,u2,v1,v2)&&!same_side(v1,v2,u1,u2);
dot_online_in(u1,v1,v2)||dot_online_in(u2,v1,v2)||dot_online_in(v1,u1,u2)||dot_online_in
(v2,u1,u2);
}
//判两线段相交,不包括端点和部分重合
int intersect_ex(line3 u,line3 v){
    return
dots\_onplane(u.a,u.b,v.a,v.b) \& opposite\_side(u.a,u.b,v) \& opposite\_side(v.a,v.b,u);
int intersect_ex(point3 u1,point3 u2,point3 v1,point3 v2){
```

```
return
dots onplane(u1,u2,v1,v2)&&opposite side(u1,u2,v1,v2)&&opposite side(v1,v2,u1,u2);
}
//判线段与空间三角形相交,包括交于边界和(部分)包含
int intersect in(line3 l,plane3 s){
   return !same_side(l.a,l.b,s)&&!same_side(s.a,s.b,l.a,l.b,s.c)&&
       !same_side(s.b,s.c,l.a,l.b,s.a)&&!same_side(s.c,s.a,l.a,l.b,s.b);
}
int intersect_in(point3 l1,point3 l2,point3 s1,point3 s2,point3 s3){
   return !same_side(l1,l2,s1,s2,s3)&&!same_side(s1,s2,l1,l2,s3)&&
       !same_side(s2,s3,l1,l2,s1)&&!same_side(s3,s1,l1,l2,s2);
}
//判线段与空间三角形相交,不包括交于边界和(部分)包含
int intersect_ex(line3 l,plane3 s){
   return opposite_side(l.a,l.b,s)&&opposite_side(s.a,s.b,l.a,l.b,s.c)&&
       opposite side(s.b,s.c,l.a,l.b,s.a)&&opposite side(s.c,s.a,l.a,l.b,s.b);
}
int intersect ex(point3 l1,point3 l2,point3 s1,point3 s2,point3 s3){
   return opposite_side(l1,l2,s1,s2,s3)&&opposite_side(s1,s2,l1,l2,s3)&&
       opposite_side(s2,s3,l1,l2,s1)&&opposite_side(s3,s1,l1,l2,s2);
}
//计算两直线交点,注意事先判断直线是否共面和平行!
//线段交点请另外判线段相交(同时还是要判断是否平行!)
point3 intersection(line3 u,line3 v){
   point3 ret=u.a;
   double t=((u.a.x-v.a.x)*(v.a.y-v.b.y)-(u.a.y-v.a.y)*(v.a.x-v.b.x))
           /((u.a.x-u.b.x)*(v.a.y-v.b.y)-(u.a.y-u.b.y)*(v.a.x-v.b.x));
   ret.x+=(u.b.x-u.a.x)*t;
   ret.y+=(u.b.y-u.a.y)*t;
   ret.z+=(u.b.z-u.a.z)*t;
   return ret;
}
point3 intersection(point3 u1,point3 u2,point3 v1,point3 v2){
   point3 ret=u1;
   double t=((u1.x-v1.x)*(v1.y-v2.y)-(u1.y-v1.y)*(v1.x-v2.x))
           /((u1.x-u2.x)*(v1.y-v2.y)-(u1.y-u2.y)*(v1.x-v2.x));
   ret.x+=(u2.x-u1.x)*t;
   ret.y+=(u2.y-u1.y)*t;
   ret.z+=(u2.z-u1.z)*t;
   return ret;
}
```

```
//线段和空间三角形交点请另外判断
point3 intersection(line3 l,plane3 s){
   point3 ret=pvec(s);
   double t=(ret.x*(s.a.x-l.a.x)+ret.y*(s.a.y-l.a.y)+ret.z*(s.a.z-l.a.z))/
       (ret.x*(l.b.x-l.a.x)+ret.y*(l.b.y-l.a.y)+ret.z*(l.b.z-l.a.z));
   ret.x=l.a.x+(l.b.x-l.a.x)*t;
   ret.y=l.a.y+(l.b.y-l.a.y)*t;
   ret.z=l.a.z+(l.b.z-l.a.z)*t;
   return ret;
}
point3 intersection(point3 l1,point3 l2,point3 s1,point3 s2,point3 s3){
   point3 ret=pvec(s1,s2,s3);
   double t=(ret.x*(s1.x-l1.x)+ret.y*(s1.y-l1.y)+ret.z*(s1.z-l1.z))/
       (ret.x*(l2.x-l1.x)+ret.y*(l2.y-l1.y)+ret.z*(l2.z-l1.z));
   ret.x=l1.x+(l2.x-l1.x)*t;
   ret.y=l1.y+(l2.y-l1.y)*t;
   ret.z=l1.z+(l2.z-l1.z)*t;
   return ret;
}
//计算两平面交线,注意事先判断是否平行,并保证三点不共线!
line3 intersection(plane3 u,plane3 v){
   line3 ret;
   ret.a=parallel(v.a,v.b,u.a,u.b,u.c)?intersection(v.b,v.c,u.a,u.b,u.c):intersection(v
.a,v.b,u.a,u.b,u.c);
   ret.b=parallel(v.c,v.a,u.a,u.b,u.c)?intersection(v.b,v.c,u.a,u.b,u.c):intersection(v
.c,v.a,u.a,u.b,u.c);
   return ret;
line3 intersection(point3 u1,point3 u2,point3 u3,point3 v1,point3 v2,point3 v3){
   line3 ret;
   ret.a=parallel(v1,v2,u1,u2,u3)?intersection(v2,v3,u1,u2,u3):intersection(v1,v2,u1,u2
,u3);
   ret.b=parallel(v3,v1,u1,u2,u3)?intersection(v2,v3,u1,u2,u3):intersection(v3,v1,u1,u2
,u3);
   return ret;
}
//点到直线距离
double ptoline(point3 p,line3 l){
   return vlen(xmult(subt(p,l.a),subt(l.b,l.a)))/distance(l.a,l.b);
}
double ptoline(point3 p,point3 l1,point3 l2){
   return vlen(xmult(subt(p,l1),subt(l2,l1)))/distance(l1,l2);
}
```

```
//点到平面距离
double ptoplane(point3 p,plane3 s){
   return fabs(dmult(pvec(s),subt(p,s.a)))/vlen(pvec(s));
}
double ptoplane(point3 p,point3 s1,point3 s2,point3 s3){
   return fabs(dmult(pvec(s1,s2,s3),subt(p,s1)))/vlen(pvec(s1,s2,s3));
}
//直线到直线距离
double linetoline(line3 u,line3 v){
   point3 n=xmult(subt(u.a,u.b),subt(v.a,v.b));
   return fabs(dmult(subt(u.a,v.a),n))/vlen(n);
}
double linetoline(point3 u1,point3 u2,point3 v1,point3 v2){
   point3 n=xmult(subt(u1,u2),subt(v1,v2));
   return fabs(dmult(subt(u1,v1),n))/vlen(n);
}
//两直线夹角 cos 值
double angle_cos(line3 u,line3 v){
   return dmult(subt(u.a,u.b),subt(v.a,v.b))/vlen(subt(u.a,u.b))/vlen(subt(v.a,v.b));
}
double angle_cos(point3 u1,point3 u2,point3 v1,point3 v2){
   return dmult(subt(u1,u2),subt(v1,v2))/vlen(subt(u1,u2))/vlen(subt(v1,v2));
}
//两平面夹角 cos 值
double angle_cos(plane3 u,plane3 v){
   return dmult(pvec(u),pvec(v))/vlen(pvec(u))/vlen(pvec(v));
}
double angle_cos(point3 u1,point3 u2,point3 u3,point3 v1,point3 v2,point3 v3){
dmult(pvec(u1,u2,u3),pvec(v1,v2,v3))/vlen(pvec(u1,u2,u3))/vlen(pvec(v1,v2,v3));
}
//直线平面夹角 sin 值
double angle_sin(line3 l,plane3 s){
   return dmult(subt(l.a,l.b),pvec(s))/vlen(subt(l.a,l.b))/vlen(pvec(s));
double angle_sin(point3 l1,point3 l2,point3 s1,point3 s2,point3 s3){
   return dmult(subt(l1,l2),pvec(s1,s2,s3))/vlen(subt(l1,l2))/vlen(pvec(s1,s2,s3));
}
                                  9. 凸包(graham)
// CONVEX HULL I
// modified by rr 不能去掉点集中重合的点
```

```
#include <stdlib.h>
#define eps 1e-8
#define zero(x) (((x)>0?(x):-(x))<eps)
struct point{double x,y;};
//计算 cross product (P1-P0)x(P2-P0)
double xmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
//graham 算法顺时针构造包含所有共线点的凸包,0(nlogn)
point p1,p2;
int graham_cp(const void* a,const void* b){
   double ret=xmult(*((point*)a),*((point*)b),p1);
    return zero(ret)?(xmult(*((point*)a),*((point*)b),p2)>0?1:-1):(ret>0?1:-1);
}
void _graham(int n,point* p,int& s,point* ch){
   int i,k=0;
   for (p1=p2=p[0], i=1; i< n; p2.x+=p[i].x, p2.y+=p[i].y, i++)
       if (p1.y-p[i].y>eps||(zero(p1.y-p[i].y)&&p1.x>p[i].x))
           p1=p[k=i];
   p2.x/=n,p2.y/=n;
   p[k]=p[0],p[0]=p1;
   qsort(p+1,n-1,sizeof(point),graham_cp);
   for (ch[0]=p[0], ch[1]=p[1], ch[2]=p[2], s=i=3; i<n; ch[s++]=p[i++])
       for (;s>2&&xmult(ch[s-2],p[i],ch[s-1])<-eps;s--);</pre>
}
//构造凸包接口函数,传入原始点集大小 n,点集 p(p 原有顺序被打乱!)
//返回凸包大小,凸包的点在 convex 中
//参数 maxsize 为 1 包含共线点,为 0 不包含共线点,缺省为 1
//参数 clockwise 为 1 顺时针构造,为 0 逆时针构造,缺省为 1
//在输入仅有若干共线点时算法不稳定,可能有此类情况请另行处理!
//不能去掉点集中重合的点
int graham(int n,point* p,point* convex,int maxsize=1,int dir=1){
   point* temp=new point[n];
   int s,i;
   _graham(n,p,s,temp);
   for (convex[0]=temp[0], n=1, i=(dir?1:(s-1)); dir?(i<s):i; i+=(dir?1:-1))
       if (maxsize||!zero(xmult(temp[i-1],temp[i],temp[(i+1)%s])))
           convex[n++]=temp[i];
   delete []temp;
   return n;
}
// CONVEX HULL II
// modified by mgmg 去掉点集中重合的点
```

```
#define eps 1e-8
#define zero(x) (((x)>0?(x):-(x))<eps)
struct point{double x,y;};
//计算 cross product (P1-P0)x(P2-P0)
double xmult(point p1,point p2,point p0){
   return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
//graham 算法顺时针构造包含所有共线点的凸包,0(nlogn)
point p1,p2;
int graham_cp(const void* a,const void* b){
   double ret=xmult(*((point*)a),*((point*)b),p1);
   return zero(ret)?(xmult(*((point*)a),*((point*)b),p2)>0?1:-1):(ret>0?1:-1);
}
void _graham(int n,point* p,int& s,point* ch){
   int i,k=0;
   for (p1=p2=p[0], i=1; i< n; p2.x+=p[i].x, p2.y+=p[i].y, i++)
      if (p1.y-p[i].y>eps||(zero(p1.y-p[i].y)&&p1.x>p[i].x))
          p1=p[k=i];
   p2.x/=n,p2.y/=n;
   p[k]=p[0],p[0]=p1;
   qsort(p+1,n-1,sizeof(point),graham_cp);
   for (ch[0]=p[0], ch[1]=p[1], ch[2]=p[2], s=i=3; i<n; ch[s++]=p[i++])
      for (;s>2&&xmult(ch[s-2],p[i],ch[s-1])<-eps;s--);</pre>
}
int wipesame_cp(const void *a, const void *b)
   if ((*(point *)a).y < (*(point *)b).y - eps) return -1;
   else if ((*(point *)a).y > (*(point *)b).y + eps) return 1;
   else if ((*(point *)a).x < (*(point *)b).x - eps) return -1;
   else if ((*(point *)a).x > (*(point *)b).x + eps) return 1;
   else return 0;
}
int _wipesame(point * p, int n)
   int i, k;
   qsort(p, n, sizeof(point), wipesame_cp);
   for (k=i=1;i<n;i++)
      if (wipesame_cp(p+i,p+i-1)!=0) p[k++]=p[i];
   return k;
}
//构造凸包接口函数,传入原始点集大小n,点集p(p原有顺序被打乱!)
```

```
//返回凸包大小,凸包的点在 convex 中
//参数 maxsize 为 1 包含共线点,为 0 不包含共线点,缺省为 1
//参数 clockwise 为 1 顺时针构造,为 0 逆时针构造,缺省为 1
//在输入仅有若干共线点时算法不稳定,可能有此类情况请另行处理!
int graham(int n,point* p,point* convex,int maxsize=1,int dir=1){
   point* temp=new point[n];
   int s,i;
   n = \_wipesame(p,n);
   _graham(n,p,s,temp);
   for (convex[0]=temp[0], n=1, i=(dir?1:(s-1)); dir?(i<s):i; i+=(dir?1:-1))
      if (maxsize||!zero(xmult(temp[i-1],temp[i],temp[(i+1)%s])))
         convex[n++]=temp[i];
   delete []temp;
   return n;
}
                                  10. 网格(pick)
#define abs(x) ((x)>0?(x):-(x))
struct point{int x,y;};
int gcd(int a,int b){
   return b?gcd(b,a%b):a;
}
//多边形上的网格点个数
int grid onedge(int n,point* p){
   int i,ret=0;
   for (i=0;i<n;i++)
       ret+=gcd(abs(p[i].x-p[(i+1)%n].x),abs(p[i].y-p[(i+1)%n].y));
   return ret;
}
//多边形内的网格点个数
int grid_inside(int n,point* p){
   int i,ret=0;
   for (i=0;i<n;i++)
       ret+=p[(i+1)%n].y*(p[i].x-p[(i+2)%n].x);
   return (abs(ret)-grid_onedge(n,p))/2+1;
}
                                       11. 圆
#include <math.h>
#define eps 1e-8
struct point{double x,y;};
double xmult(point p1,point p2,point p0){
```

```
return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
double distance(point p1,point p2){
    return sqrt((p1.x-p2.x)*(p1.x-p2.x)+(p1.y-p2.y)*(p1.y-p2.y));
}
double disptoline(point p,point l1,point l2){
    return fabs(xmult(p,l1,l2))/distance(l1,l2);
}
point intersection(point u1,point u2,point v1,point v2){
    point ret=u1;
    double t=((u1.x-v1.x)*(v1.y-v2.y)-(u1.y-v1.y)*(v1.x-v2.x))
           /((u1.x-u2.x)*(v1.y-v2.y)-(u1.y-u2.y)*(v1.x-v2.x));
    ret.x+=(u2.x-u1.x)*t;
    ret.y+=(u2.y-u1.y)*t;
    return ret;
}
//判直线和圆相交,包括相切
int intersect_line_circle(point c,double r,point l1,point l2){
    return disptoline(c,l1,l2)<r+eps;</pre>
}
//判线段和圆相交,包括端点和相切
int intersect seg circle(point c,double r,point l1,point l2){
    double t1=distance(c,l1)-r,t2=distance(c,l2)-r;
    point t=c;
    if (t1<eps||t2<eps)</pre>
        return t1>-eps||t2>-eps;
    t.x+=l1.y-l2.y;
    t.y+=l2.x-l1.x;
    return xmult(l1,c,t)*xmult(l2,c,t)<eps&&disptoline(c,l1,l2)-r<eps;</pre>
}
//判圆和圆相交,包括相切
int intersect_circle_circle(point c1,double r1,point c2,double r2){
    return distance(c1,c2)<r1+r2+eps&&distance(c1,c2)>fabs(r1-r2)-eps;
}
//计算圆上到点 p 最近点,如 p 与圆心重合,返回 p 本身
point dot to circle(point c,double r,point p){
    point u,v;
    if (distance(p,c)<eps)</pre>
        return p;
```

```
u.x=c.x+r*fabs(c.x-p.x)/distance(c,p);
   u.y=c.y+r*fabs(c.y-p.y)/distance(c,p)*((c.x-p.x)*(c.y-p.y)<0?-1:1);
   v.x=c.x-r*fabs(c.x-p.x)/distance(c,p);
   v.y=c.y-r*fabs(c.y-p.y)/distance(c,p)*((c.x-p.x)*(c.y-p.y)<0?-1:1);
   return distance(u,p)<distance(v,p)?u:v;</pre>
}
//计算直线与圆的交点,保证直线与圆有交点
//计算线段与圆的交点可用这个函数后判点是否在线段上
void intersection_line_circle(point c,double r,point l1,point l2,point& p1,point& p2){
   point p=c;
   double t;
   p.x+=l1.y-l2.y;
   p.y+=l2.x-l1.x;
   p=intersection(p,c,l1,l2);
   t=sqrt(r*r-distance(p,c)*distance(p,c))/distance(l1,l2);
   p1.x=p.x+(l2.x-l1.x)*t;
   p1.y=p.y+(l2.y-l1.y)*t;
   p2.x=p.x-(l2.x-l1.x)*t;
   p2.y=p.y-(l2.y-l1.y)*t;
}
//计算圆与圆的交点,保证圆与圆有交点,圆心不重合
void intersection_circle_circle(point c1,double r1,point c2,double r2,point& p1,point& p2){
   point u,v;
   double t;
   t=(1+(r1*r1-r2*r2)/distance(c1,c2)/distance(c1,c2))/2;
   u.x=c1.x+(c2.x-c1.x)*t;
   u.y=c1.y+(c2.y-c1.y)*t;
   v.x=u.x+c1.y-c2.y;
   v.y=u.y-c1.x+c2.x;
   intersection_line_circle(c1,r1,u,v,p1,p2);
}
                                    12. 整数函数
```

```
//整数几何函数库
//注意某些情况下整数运算会出界!
#define sign(a) ((a)>0?1:(((a)<0?-1:0)))
struct point{int x,y;};
struct line{point a,b;};

//计算 cross product (P1-P0)x(P2-P0)
int xmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.y-p0.y)-(p2.x-p0.x)*(p1.y-p0.y);
}
```

```
int xmult(int x1,int y1,int x2,int y2,int x0,int y0){
    return (x1-x0)*(y2-y0)-(x2-x0)*(y1-y0);
}
//计算 dot product (P1-P0).(P2-P0)
int dmult(point p1,point p2,point p0){
    return (p1.x-p0.x)*(p2.x-p0.x)+(p1.y-p0.y)*(p2.y-p0.y);
}
int dmult(int x1,int y1,int x2,int y2,int x0,int y0){
    return (x1-x0)*(x2-x0)+(y1-y0)*(y2-y0);
}
//判三点共线
int dots_inline(point p1,point p2,point p3){
    return !xmult(p1,p2,p3);
int dots_inline(int x1,int y1,int x2,int y2,int x3,int y3){
    return !xmult(x1,y1,x2,y2,x3,y3);
}
//判点是否在线段上,包括端点和部分重合
int dot_online_in(point p,line l){
    \label{eq:continuity} return \ !xmult(p,l.a,l.b)\&\&(l.a.x-p.x)*(l.b.x-p.x)<=0\&\&(l.a.y-p.y)*(l.b.y-p.y)<=0;
}
int dot_online_in(point p,point l1,point l2){
    return \ !xmult(p,l1,l2)\&\&(l1.x-p.x)*(l2.x-p.x)<=0\&\&(l1.y-p.y)*(l2.y-p.y)<=0;
}
int dot_online_in(int x,int y,int x1,int y1,int x2,int y2){
    return !xmult(x,y,x1,y1,x2,y2)&&(x1-x)*(x2-x)<=0&&(y1-y)*(y2-y)<=0;
}
//判点是否在线段上,不包括端点
int dot_online_ex(point p,line l){
    return dot_online_in(p,l)&&(p.x!=l.a.x||p.y!=l.a.y)&&(p.x!=l.b.x||p.y!=l.b.y);
}
int dot_online_ex(point p,point l1,point l2){
    return dot_online_in(p,l1,l2)&&(p.x!=l1.x||p.y!=l1.y)&&(p.x!=l2.x||p.y!=l2.y);
int dot_online_ex(int x,int y,int x1,int y1,int x2,int y2){
    return dot_online_in(x,y,x1,y1,x2,y2)&&(x!=x1||y!=y1)&&(x!=x2||y!=y2);
}
//判两点在直线同侧,点在直线上返回 0
int same_side(point p1,point p2,line l){
    return sign(xmult(l.a,p1,l.b))*xmult(l.a,p2,l.b)>0;
}
```

```
int same_side(point p1,point p2,point l1,point l2){
          return sign(xmult(l1,p1,l2))*xmult(l1,p2,l2)>0;
}
//判两点在直线异侧,点在直线上返回 0
int opposite_side(point p1,point p2,line l){
          return sign(xmult(l.a,p1,l.b))*xmult(l.a,p2,l.b)<0;</pre>
}
int opposite side(point p1,point p2,point l1,point l2){
          return sign(xmult(l1,p1,l2))*xmult(l1,p2,l2)<0;</pre>
}
//判两直线平行
int parallel(line u,line v){
          return (u.a.x-u.b.x)*(v.a.y-v.b.y)==(v.a.x-v.b.x)*(u.a.y-u.b.y);
}
int parallel(point u1,point u2,point v1,point v2){
          return (u1.x-u2.x)*(v1.y-v2.y)==(v1.x-v2.x)*(u1.y-u2.y);
}
//判两直线垂直
int perpendicular(line u,line v){
          return (u.a.x-u.b.x)*(v.a.x-v.b.x)==-(u.a.y-u.b.y)*(v.a.y-v.b.y);
}
int perpendicular(point u1,point u2,point v1,point v2){
          return (u1.x-u2.x)*(v1.x-v2.x)==-(u1.y-u2.y)*(v1.y-v2.y);
}
//判两线段相交,包括端点和部分重合
int intersect_in(line u,line v){
         if (!dots_inline(u.a,u.b,v.a)||!dots_inline(u.a,u.b,v.b))
                    return !same_side(u.a,u.b,v)&&!same_side(v.a,v.b,u);
          return
dot_online_in(u.a,v)||dot_online_in(u.b,v)||dot_online_in(v.a,u)||dot_online_in(v.b,u);
int intersect_in(point u1,point u2,point v1,point v2){
         if (!dots_inline(u1,u2,v1)||!dots_inline(u1,u2,v2))
                    return !same_side(u1,u2,v1,v2)&&!same_side(v1,v2,u1,u2);
\label{localine_in} dot\_online\_in(u1,v1,v2) | | dot\_online\_in(u2,v1,v2) | | dot\_online\_in(v1,u1,u2) | | dot\_online\_in(u2,v1,v2) | | dot\_online\_in(u2,v1,v2) | | dot\_online\_in(u2,v1,v2) | | dot\_online\_in(v1,u1,u2) | | dot\_online\_in(u2,v1,v2) | | dot\_online\_in(u2,v1,v2) | | dot\_online\_in(v1,u1,u2) | | dot\_online\_in(u2,v1,v2) | | dot\_onli
(v2,u1,u2);
}
//判两线段相交,不包括端点和部分重合
int intersect_ex(line u,line v){
```

```
return opposite_side(u.a,u.b,v)&&opposite_side(v.a,v.b,u);
}
int intersect_ex(point u1,point u2,point v1,point v2){
    return opposite_side(u1,u2,v1,v2)&&opposite_side(v1,v2,u1,u2);
}
```

13. 注意

- 1. 注意舍入方式(0.5 的舍入方向);防止输出-0.
- 2. 几何题注意多测试不对称数据.
- 3. 整数几何注意 xmult 和 dmult 是否会出界; 符点几何注意 eps 的使用.
- 4. 避免使用斜率;注意除数是否会为 0.
- 5. 公式一定要化简后再代入.
- 6. 判断同一个 2*PI 域内两角度差应该是 abs(a1-a2)
beta||abs(a1-a2)>pi+pi-beta; 相等应该是 abs(a1-a2)<eps||abs(a1-a2)>pi+pi-eps;
- 7. 需要的话尽量使用 atan2,注意:atan2(0,0)=0, atan2(1,0)=pi/2,atan2(-1,0)=-pi/2,atan2(0,1)=0,atan2(0,-1)=pi.
- 8. cross product = |u|*|v|*sin(a)dot product = |u|*|v|*cos(a)
- 9. (P1-P0)x(P2-P0)结果的意义:

正: <P0,P1>在<P0,P2>顺时针(0,pi)内 负: <P0,P1>在<P0,P2>逆时针(0,pi)内 0: <P0,P1>,<P0,P2>共线,夹角为0或pi

10. 误差限缺省使用 1e-8!

十二。结构

1. 并查集

//带路径压缩的并查集,用于动态维护查询等价类
//图论算法中动态判点集连通常用
//维护和查询复杂度略大于 0(1)
//集合元素取值 1..MAXN-1(注意 0 不能用!),默认不等价
#include <string.h>

```
#define MAXN 100000
#define _ufind_run(x) for(;p[t=x];x=p[x],p[t]=(p[x]?p[x]:x))
#define _run_both _ufind_run(i);_ufind_run(j)
struct ufind{
   int p[MAXN],t;
   void init(){memset(p,0,sizeof(p));}
   void set_friend(int i,int j){_run_both;p[i]=(i==j?0:j);}
   int is_friend(int i,int j){_run_both;return i==j&&i;}
};
                          2. 并查集扩展(friend_enemy)
//带路径压缩的并查集扩展形式
//用于动态维护查询 friend-enemy 型等价类
//维护和查询复杂度略大于 0(1)
//集合元素取值 1..MAXN-1(注意 0 不能用!),默认无关
#include <string.h>
#define MAXN 100000
#define sig(x) ((x)>0?1:-1)
#define abs(x) ((x)>0?(x):-(x))
#define
                                                                        _ufind_run(x)
for(;p[t=abs(x)];x=sig(x)*p[abs(x)],p[t]=sig(p[t])*(p[abs(x)]?p[abs(x)]:abs(p[t])))
#define _run_both _ufind_run(i);_ufind_run(j)
#define _set_side(x) p[abs(i)]=sig(i)*(abs(i)==abs(j)?0:(x)*j)
#define judge side(x) (i==(x)*j\&\&i)
struct ufind{
   int p[MAXN],t;
   void init(){memset(p,0,sizeof(p));}
   int set_friend(int i,int j){_run_both;_set_side(1);return !_judge_side(-1);}
   int set_enemy(int i,int j){_run_both;_set_side(-1);return !_judge_side(1);}
   int is_friend(int i,int j){_run_both;return _judge_side(1);}
   int is_enemy(int i,int j){_run_both;return _judge_side(-1);}
};
                                  3. 堆(binary)
//二分堆(binary)
//可插入,获取并删除最小(最大)元素,复杂度均 0(logn)
//可更改元素类型,修改比较符号或换成比较函数
#define MAXN 10000
#define _{cp(a,b)}((a)<(b))
typedef int elem_t;
struct heap{
   elem_t h[MAXN];
```

```
int n,p,c;
          void init(){n=0;}
          void ins(elem t e){
                     for (p=++n;p>1&\&_cp(e,h[p>>1]);h[p]=h[p>>1],p>>=1);
                     h[p]=e;
          }
          int del(elem_t& e){
                     if (!n) return 0;
(e=h[p=1], c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]);h[p]=h[c],p=c,c<<=1);
                     h[p]=h[n--];
                     return 1;
          }
};
                                                                                                 4. 堆(mapped)
//映射二分堆
//可插入,获取并删除任意元素,复杂度均 0(logn)
//插入时提供一个索引值,删除时按该索引删除,获取并删除最小元素时一起获得该索引
//索引值范围 0..MAXN-1,不能重复,不负责维护索引的唯一性,不在此返回请另外映射
//主要用于图论算法,该索引值可以是节点的下标
//可更改元素类型,修改比较符号或换成比较函数
#define MAXN 10000
#define _cp(a,b) ((a)<(b))
typedef int elem_t;
struct heap{
          elem t h[MAXN];
          int ind[MAXN],map[MAXN],n,p,c;
          void init(){n=0;}
          void ins(int i,elem t e){
                     for (p=++n;p>1&\&\_cp(e,h[p>>1]);h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
                     h[map[ind[p]=i]=p]=e;
          }
          int del(int i,elem t& e){
                     i=map[i];if (i<1||i>n) return 0;
                     for (e=h[p=i];p>1;h[map[ind[p]=ind[p>>1]]=p]=h[p>>1],p>>=1);
                     for
(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]); h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]); h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]); h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]); h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]); h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+1],h[c]))],h[n]); h[map[ind[p]=ind[c]]=p]=h[c],p=c,c<n>(c=2; c<n&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c+=(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&\_cp(h[c==(c<n-1&\&
=1);
                     h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
          }
          int delmin(int& i,elem_t& e){
                     if (n<1) return 0;i=ind[1];</pre>
                     for
(e=h[p=1], c=2; c<n&_cp(h[c+=(c<n-1&&_cp(h[c+1],h[c]))],h[n]);h[map[ind[p]=ind[c]]=p]=h[c]
```

```
],p=c,c<<=1);
       h[map[ind[p]=ind[n]]=p]=h[n];n--;return 1;
   }
};
                                    5. 矩形切割
//矩形切割
//intersect 函数构造矩形 a 和 b 的交集
//cut 函数将 b 关于 a 进行切割,用切下的矩形作参数调用 dummy 函数
struct rect{
   int x1,x2,y1,y2;
   rect(){}
   rect(int a,int b,int c,int d):x1(a),x2(b),y1(c),y2(d){}
};
inline rect intersect(rect& a,rect& b){
    return
rect(a.x1>b.x1?a.x1:b.x1,a.x2<b.x2?a.x2:b.x2,a.y1>b.y1?a.y1:b.y1,a.y2<b.y2?a.y2:b.y2);
void dummy(rect a){
   //dispose this rect
}
int cut(rect& a,rect b){
   rect c=intersect(a,b);
   if (c.x1>=c.x2||c.y1>=c.y2)
       return 0;
   if (b.x1<c.x1)
       dummy(rect(b.x1,c.x1,b.y1,b.y2)),b.x1=c.x1;
   if (b.x2>c.x2)
       dummy(rect(c.x2,b.x2,b.y1,b.y2)),b.x2=c.x2;
   if (b.y1<c.y1)
       dummy(rect(b.x1,b.x2,b.y1,c.y1)),b.y1=c.y1;
   if (b.y2>c.y2)
       dummy(rect(b.x1,b.x2,c.y2,b.y2)),b.y2=c.y2;
   return 1;
}
                                     6. 线段树
//线段树
//可以处理加入边和删除边不同的情况
//inc_seg 和 dec_seg 用于加入边
//seg_len 求长度
//t 传根节点(一律为 1)
//l0,r0 传树的节点范围(一律为 1..t)
```

```
//l,r传线段(端点)
#define MAXN 10000
struct segtree{
    int n,cnt[MAXN],len[MAXN];
    segtree(int t):n(t){
        for (int i=1;i<=t;i++)</pre>
            cnt[i]=len[i]=0;
    };
    void update(int t,int l,int r);
    void inc_seg(int t,int l0,int r0,int l,int r);
    void dec_seg(int t,int l0,int r0,int l,int r);
    int seg_len(int t,int l0,int r0,int l,int r);
};
int length(int l,int r){
    return r-l;
}
void segtree::update(int t,int l,int r){
    if (cnt[t]||r-l==1)
        len[t]=length(l,r);
    else
        len[t]=len[t+t]+len[t+t+1];
}
void segtree::inc_seg(int t,int l0,int r0,int l,int r){
    if (l0==l\&\&r0==r)
        cnt[t]++;
    else{
        int m0=(l0+r0)>>1;
        if (l<m0)
            inc_seg(t+t,l0,m0,l,m0<r?m0:r);</pre>
        if (r>m0)
            inc seg(t+t+1,m0,r0,m0>l?m0:l,r);
        if (cnt[t+t]&&cnt[t+t+1]){
            cnt[t+t]--;
            update(t+t,l0,m0);
            cnt[t+t+1]--;
            update(t+t+1,m0,r0);
            cnt[t]++;
        }
    update(t,l0,r0);
}
void segtree::dec_seg(int t,int l0,int r0,int l,int r){
```

```
if (l0==l\&\&r0==r)
       cnt[t]--;
   else if (cnt[t]){
       cnt[t]--;
       if (l>l0)
           inc_seg(t,l0,r0,l0,l);
       if (r<r0)
           inc_seg(t,l0,r0,r,r0);
   }
   else{
       int m0=(l0+r0)>>1;
       if (l<m0)
           dec_seg(t+t,l0,m0,l,m0<r?m0:r);</pre>
       if (r>m0)
           dec_seg(t+t+1,m0,r0,m0>l?m0:l,r);
   }
   update(t,l0,r0);
}
int segtree::seg_len(int t,int l0,int r0,int l,int r){
   if (cnt[t]||(l0==l\&r0==r))
       return len[t];
   else{
       int m0=(l0+r0)>>1, ret=0;
       if (l<m0)
           ret+=seg_len(t+t,l0,m0,l,m0<r?m0:r);
       if (r>m0)
           ret+=seg_len(t+t+1,m0,r0,m0>l?m0:l,r);
       return ret;
   }
}
                                    7. 线段树扩展
//线段树扩展
//可以计算长度和线段数
//可以处理加入边和删除边不同的情况
//inc_seg 和 dec_seg 用于加入边
//seg_len 求长度,seg_cut 求线段数
//t 传根节点(一律为 1)
//l0,r0 传树的节点范围(一律为 1..t)
//l,r传线段(端点)
#define MAXN 10000
struct segtree{
   int n,cnt[MAXN],len[MAXN],cut[MAXN],bl[MAXN],br[MAXN];
   segtree(int t):n(t){
```

for (int i=1;i<=t;i++)

```
cnt[i]=len[i]=cut[i]=bl[i]=br[i]=0;
    };
    void update(int t,int l,int r);
    void inc_seg(int t,int l0,int r0,int l,int r);
    void dec_seg(int t,int l0,int r0,int l,int r);
    int seg_len(int t,int l0,int r0,int l,int r);
    int seg_cut(int t,int l0,int r0,int l,int r);
};
int length(int l,int r){
    return r-l;
}
void segtree::update(int t,int l,int r){
    if (cnt[t]||r-l==1)
        len[t]=length(l,r),cut[t]=bl[t]=br[t]=1;
   else{
        len[t]=len[t+t]+len[t+t+1];
        cut[t]=cut[t+t]+cut[t+t+1];
        if (br[t+t]&&bl[t+t+1])
            cut[t]--;
        bl[t]=bl[t+t],br[t]=br[t+t+1];
   }
}
void segtree::inc_seg(int t,int l0,int r0,int l,int r){
    if (l0==l\&\&r0==r)
        cnt[t]++;
    else{
        int m0=(l0+r0)>>1;
        if (l<m0)
            inc_seg(t+t,l0,m0,l,m0<r?m0:r);
        if (r>m0)
            inc seg(t+t+1,m0,r0,m0>l?m0:l,r);
        if (cnt[t+t]&&cnt[t+t+1]){
            cnt[t+t]--;
            update(t+t,l0,m0);
            cnt[t+t+1]--;
            update(t+t+1,m0,r0);
           cnt[t]++;
       }
    update(t,l0,r0);
}
void segtree::dec_seg(int t,int l0,int r0,int l,int r){
```

```
if (l0==l\&\&r0==r)
        cnt[t]--;
    else if (cnt[t]){
        cnt[t]--;
        if (l>l0)
            inc_seg(t,l0,r0,l0,l);
        if (r<r0)
            inc_seg(t,l0,r0,r,r0);
    }
   else{
        int m0=(l0+r0)>>1;
        if (l<m0)
            dec_seg(t+t,l0,m0,l,m0<r?m0:r);</pre>
        if (r>m0)
            dec_seg(t+t+1,m0,r0,m0>l?m0:l,r);
    }
    update(t,l0,r0);
}
int segtree::seg_len(int t,int l0,int r0,int l,int r){
    if (cnt[t]||(l0==l\&r0==r))
        return len[t];
    else{
        int m0=(l0+r0)>>1, ret=0;
        if (l<m0)
            ret+=seg_len(t+t,l0,m0,l,m0<r?m0:r);
        if (r>m0)
            ret+=seg_len(t+t+1,m0,r0,m0>l?m0:l,r);
        return ret;
    }
}
int segtree::seg_cut(int t,int l0,int r0,int l,int r){
    if (cnt[t])
        return 1;
    if (l0==l\&\&r0==r)
        return cut[t];
    else{
        int m0=(l0+r0)>>1, ret=0;
        if (l<m0)
            ret+=seg_cut(t+t,l0,m0,l,m0<r?m0:r);
        if (r>m0)
            ret+=seg_cut(t+t+1,m0,r0,m0>l?m0:l,r);
        if (1<m0&&r>m0&&br[t+t]&&bl[t+t+1])
            ret--;
        return ret;
```

```
}
```

8. 线段树应用

求面积:

- 1) 坐标离散化
- 2) 垂直边按 x 坐标排序
- 3) 从左往右用线段树处理垂直边 累计每个离散 x 区间长度和线段树长度的乘积

求周长:

- 1) 坐标离散化
- 2) 垂直边按 x 坐标排序, 第二关键字为入边优于出边
- 3) 从左往右用线段树处理垂直边 在每个离散点上先加入所有入边,累计线段树长度变化值 再删除所有出边,累计线段树长度变化值
- 4) 水平边按 y 坐标排序, 第二关键字为入边优于出边
- 5) 从上往下用线段树处理水平边 在每个离散点上先加入所有入边,累计线段树长度变化值 再删除所有出边,累计线段树长度变化值

9. 子段和

```
//求 sum{[0..n-1]}
//维护和查询复杂度均为 O(logn)
//用于动态求子段和,数组内容保存在 sum.a[]中
//可以改成其他数据类型
#include <string.h>
#define lowbit(x) ((x)&((x)^((x)-1)))
#define MAXN 10000
typedef int elem_t;

struct sum{
    elem_t a[MAXN],c[MAXN],ret;
    int n;
    void init(int i){memset(a,0,sizeof(a));memset(c,0,sizeof(c));n=i;}
    void update(int i,elem_t v){for (v-=a[i],a[i++]+=v;i<=n;c[i-1]+=v,i+=lowbit(i));}
    elem_t query(int i){for (ret=0;i;ret+=c[i-1],i^=lowbit(i));return ret;}
};
```

10. 子阵和

```
//求 sum{a[0..m-1][0..n-1]}
//维护和查询复杂度均为 0(logm*logn)
//用于动态求子阵和,数组内容保存在 sum.a[][]中
//可以改成其他数据类型
#include <string.h>
```

```
#define lowbit(x) ((x)&((x)^{(x)-1)})
#define MAXN 100
typedef int elem_t;
struct sum{
    elem_t a[MAXN][MAXN],c[MAXN][MAXN],ret;
    int m,n,t;
    void init(int i,int j){memset(a,0,sizeof(a));memset(c,0,sizeof(c));m=i,n=j;}
    void update(int i,int j,elem_t v){
        for (v-=a[i][j],a[i++][j++]+=v,t=j;i<=m;i+=lowbit(i))</pre>
            for (j=t; j \le n; c[i-1][j-1] += v, j+=lowbit(j));
    }
    elem_t query(int i,int j){
        for (ret=0,t=j;i;i^=lowbit(i))
            for (j=t;j;ret+=c[i-1][j-1],j^=lowbit(j));
        return ret;
   }
};
```

十三. 其他

1. 分数

```
struct frac{
    int num,den;
};
double fabs(double x){
    return x>0?x:-x;
}
int gcd(int a,int b){
    int t;
    if (a<0)
       a=-a;
    if (b<0)
        b=-b;
    if (!b)
        return a;
   while (t=a%b)
        a=b,b=t;
    return b;
}
void simplify(frac& f){
```

```
int t;
    if (t=gcd(f.num,f.den))
        f.num/=t,f.den/=t;
   else
        f.den=1;
}
frac f(int n,int d,int s=1){
    frac ret;
    if (d<0)
        ret.num=-n,ret.den=-d;
    else
        ret.num=n,ret.den=d;
    if (s)
        simplify(ret);
    return ret;
}
frac convert(double x){
    frac ret;
    for (ret.den=1; fabs(x-int(x))>le-10; ret.den*=10, x*=10);
    ret.num=(int)x;
    simplify(ret);
    return ret;
}
int fraqcmp(frac a,frac b){
    int \ g1=gcd(a.den,b.den)\,,g2=gcd(a.num,b.num)\,;\\
    if (!g1||!g2)
        return 0;
    return b.den/g1*(a.num/g2)-a.den/g1*(b.num/g2);
}
frac add(frac a,frac b){
    int g1=gcd(a.den,b.den),g2,t;
        return f(1,0,0);
    t=b.den/g1*a.num+a.den/g1*b.num;
    g2=gcd(g1,t);
    return f(t/g2,a.den/g1*(b.den/g2),0);
}
frac sub(frac a,frac b){
    return add(a,f(-b.num,b.den,0));
}
```

```
frac mul(frac a,frac b){
    int t1=gcd(a.den,b.num),t2=gcd(a.num,b.den);
    if (!t1||!t2)
        return f(1,1,0);
    return f(a.num/t2*(b.num/t1),a.den/t1*(b.den/t2),0);
}
frac div(frac a,frac b){
    return mul(a,f(b.den,b.num,0));
}
                                         2. 矩阵
define MAXN 100
#define fabs(x) ((x)>0?(x):-(x))
#define zero(x) (fabs(x)<1e-10)
struct mat{
    int n,m;
    double data[MAXN][MAXN];
};
int mul(mat& c,const mat& a,const mat& b){
    int i,j,k;
    if (a.m!=b.n)
        return 0;
    c.n=a.n,c.m=b.m;
    for (i=0; i< c.n; i++)
        for (j=0; j< c.m; j++)
            for (c.data[i][j]=k=0;k<a.m;k++)
                c.data[i][j]+=a.data[i][k]*b.data[k][j];
    return 1;
}
int inv(mat& a){
    int i,j,k,is[MAXN],js[MAXN];
    double t;
    if (a.n!=a.m)
        return 0;
    for (k=0; k<a.n; k++) {
        for (t=0,i=k;i<a.n;i++)
            for (j=k;j<a.n;j++)
                if (fabs(a.data[i][j])>t)
                    t=fabs(a.data[is[k]=i][js[k]=j]);
        if (zero(t))
            return 0;
```

```
if (is[k]!=k)
            for (j=0; j<a.n; j++)
                t=a.data[k][j],a.data[k][j]=a.data[is[k]][j],a.data[is[k]][j]=t;
        if (js[k]!=k)
            for (i=0;i<a.n;i++)</pre>
                t=a.data[i][k],a.data[i][k]=a.data[i][js[k]],a.data[i][js[k]]=t;
        a.data[k][k]=1/a.data[k][k];
        for (j=0; j<a.n; j++)
            if (j!=k)
                a.data[k][j]*=a.data[k][k];
        for (i=0;i<a.n;i++)
            if (i!=k)
                for (j=0; j<a.n; j++)
                    if (j!=k)
                        a.data[i][j]-=a.data[i][k]*a.data[k][j];
        for (i=0;i<a.n;i++)
            if (i!=k)
                a.data[i][k]*=-a.data[k][k];
    }
    for (k=a.n-1;k>=0;k--){
        for (j=0; j<a.n; j++)
            if (js[k]!=k)
                t=a.data[k][j],a.data[k][j]=a.data[js[k]][j],a.data[js[k]][j]=t;
        for (i=0;i<a.n;i++)
            if (is[k]!=k)
                t=a.data[i][k],a.data[i][k]=a.data[i][is[k]],a.data[i][is[k]]=t;
    }
    return 1;
}
double det(const mat& a){
    int i,j,k,sign=0;
    double b[MAXN][MAXN],ret=1,t;
    if (a.n!=a.m)
        return 0;
    for (i=0;i<a.n;i++)
        for (j=0; j<a.m; j++)
            b[i][j]=a.data[i][j];
    for (i=0;i<a.n;i++){
        if (zero(b[i][i])){
            for (j=i+1;j<a.n;j++)
                if (!zero(b[j][i]))
                    break;
            if (j==a.n)
                return 0;
            for (k=i; k<a.n; k++)
```

```
t=b[i][k],b[i][k]=b[j][k],b[j][k]=t;
            sign++;
        }
        ret*=b[i][i];
        for (k=i+1; k<a.n; k++)
            b[i][k]/=b[i][i];
        for (j=i+1;j<a.n;j++)
           for (k=i+1; k<a.n; k++)
               b[j][k]-=b[j][i]*b[i][k];
   }
    if (sign&1)
        ret=-ret;
    return ret;
}
                                         3. 日期
//日期函数
int days[12]={31,28,31,30,31,30,31,30,31,30,31};
struct date{
   int year, month, day;
};
//判闰年
inline int leap(int year){
    return (year%4==0&&year%100!=0)||year%400==0;
}
//判合法性
inline int legal(date a){
    if (a.month<0||a.month>12)
        return 0;
   if (a.month==2)
        return a.day>0&&a.day<=28+leap(a.year);</pre>
    return a.day>0&&a.day<=days[a.month-1];</pre>
}
//比较日期大小
inline int datecmp(date a,date b){
  if (a.year!=b.year)
      return a.year-b.year;
  if (a.month!=b.month)
      return a.month-b.month;
  return a.day-b.day;
}
```

```
//返回指定日期是星期几
int weekday(date a){
    int tm=a.month>=3?(a.month-2):(a.month+10);
    int ty=a.month>=3?a.year:(a.year-1);
    return (ty+ty/4-ty/100+ty/400+(int)(2.6*tm-0.2)+a.day)%7;
}
//日期转天数偏移
int date2int(date a){
    int ret=a.year*365+(a.year-1)/4-(a.year-1)/100+(a.year-1)/400,i;
    days[1]+=leap(a.year);
    for (i=0;i<a.month-1;ret+=days[i++]);</pre>
   days[1]=28;
    return ret+a.day;
}
//天数偏移转日期
date int2date(int a){
    date ret;
    ret.year=a/146097*400;
    for (a%=146097;a>=365+leap(ret.year);a-=365+leap(ret.year),ret.year++);
    days[1]+=leap(ret.year);
    for (ret.month=1;a>=days[ret.month-1];a-=days[ret.month-1],ret.month++);
    days[1]=28;
    ret.day=a+1;
    return ret;
}
                                4. 线性方程组(gauss)
#define MAXN 100
#define fabs(x) ((x)>0?(x):-(x))
#define eps 1e-10
//列主元 gauss 消去求解 a[][]x[]=b[]
//返回是否有唯一解,若有解在 b[]中
int gauss_cpivot(int n,double a[][MAXN],double b[]){
    int i,j,k,row;
    double maxp,t;
    for (k=0; k< n; k++) {
       for (maxp=0, i=k; i< n; i++)
            if (fabs(a[i][k])>fabs(maxp))
               maxp=a[row=i][k];
       if (fabs(maxp)<eps)</pre>
            return 0;
       if (row!=k){
           for (j=k;j<n;j++)
```

```
t=a[k][j],a[k][j]=a[row][j],a[row][j]=t;
            t=b[k],b[k]=b[row],b[row]=t;
        }
        for (j=k+1; j< n; j++) {
            a[k][j]/=maxp;
            for (i=k+1;i<n;i++)</pre>
                a[i][j]-=a[i][k]*a[k][j];
        }
        b[k]/=maxp;
        for (i=k+1;i<n;i++)
            b[i]-=b[k]*a[i][k];
    }
    for (i=n-1;i>=0;i--)
        for (j=i+1;j<n;j++)
            b[i]-=a[i][j]*b[j];
    return 1;
}
//全主元 gauss 消去解 a[][]x[]=b[]
//返回是否有唯一解,若有解在 b[]中
int gauss_tpivot(int n,double a[][MAXN],double b[]){
    int i,j,k,row,col,index[MAXN];
    double maxp,t;
    for (i=0;i<n;i++)
        index[i]=i;
    for (k=0; k< n; k++) {
       for (maxp=0, i=k; i< n; i++)
            for (j=k;j<n;j++)
                if (fabs(a[i][j])>fabs(maxp))
                    maxp=a[row=i][col=j];
       if (fabs(maxp)<eps)</pre>
            return 0;
        if (col!=k){
            for (i=0;i<n;i++)
                t=a[i][col],a[i][col]=a[i][k],a[i][k]=t;
            j=index[col],index[col]=index[k],index[k]=j;
        }
        if (row!=k){
            for (j=k;j<n;j++)
             t=a[k][j],a[k][j]=a[row][j],a[row][j]=t;
            t=b[k],b[k]=b[row],b[row]=t;
        for (j=k+1; j< n; j++) {
            a[k][j]/=maxp;
            for (i=k+1;i<n;i++)
                a[i][j]-=a[i][k]*a[k][j];
```

```
}
        b[k]/=maxp;
        for (i=k+1;i<n;i++)</pre>
            b[i]-=b[k]*a[i][k];
    }
    for (i=n-1;i>=0;i--)
        for (j=i+1;j<n;j++)
            b[i]-=a[i][j]*b[j];
    for (k=0; k< n; k++)
        a[0][index[k]]=b[k];
    for (k=0; k< n; k++)
        b[k]=a[0][k];
    return 1;
}
                                        5. 线性相关
//判线性相关(正交化)
//传入m个n维向量
#include <math.h>
#define MAXN 100
#define eps 1e-10
int linear_dependent(int m,int n,double vec[][MAXN]){
    double ort[MAXN][MAXN],e;
    int i,j,k;
    if (m>n)
        return 1;
    for (i=0; i< m; i++){
        for (j=0; j< n; j++)
            ort[i][j]=vec[i][j];
        for (k=0; k<i; k++){
            for (e=j=0; j< n; j++)
                e+=ort[i][j]*ort[k][j];
            for (j=0; j< n; j++)
                ort[i][j]-=e*ort[k][j];
            for (e=j=0; j< n; j++)
                e+=ort[i][j]*ort[i][j];
            if (fabs(e=sqrt(e))<eps)</pre>
                return 1;
            for (j=0; j< n; j++)
                ort[i][j]/=e;
        }
    }
    return 0;
}
```

1. joseph

```
// Joseph's Problem
// input: n,m -- the number of persons, the inteval between persons
// output: -- return the reference of last person
int josephus0(int n, int m)
{
   if (n == 2) return (m\%2) ? 2 : 1;
    int v = (m+josephus0(n-1,m)) % n;
    if (v == 0) v = n;
    return v;
}
int josephus(int n, int m)
    if (m == 1) return n;
   if (n == 1) return 1;
   if (m >=n) return josephus0(n,m);
   int l = (n/m)*m;
   int j = josephus(n - (n/m), m);
    if (j <= n-l) return l+j;</pre>
    j -= n-l;
   int t = (j/(m-1))*m;
    if ((j % (m-1)) == 0) return t-1;
    return t + (j % (m-1));
}
```

2. N 皇后构造解

```
void even1(int n,int *p){
   int i;
   for (i=1;i<=n/2;i++)
       p[i-1]=2*i;
   for (i=n/2+1;i<=n;i++)
       p[i-1]=2*i-n-1;
}

void even2(int n,int *p){
   int i;
   for (i=1;i<=n/2;i++)
       p[i-1]=(2*i+n/2-3)%n+1;</pre>
```

//N 皇后构造解, n>=4

```
for (i=n/2+1; i \le n; i++)
       p[i-1]=n-(2*(n-i+1)+n/2-3)%n;
}
void generate(int,int*);
void odd(int n,int *p){
   generate(n-1,p),p[n-1]=n;
}
void generate(int n,int *p){
    if (n&1)
       odd(n,p);
    else if (n%6!=2)
       even1(n,p);
   else
       even2(n,p);
}
                                    3. 布尔母函数
//布尔母函数
//判 m[]个价值为 w[]的货币能否构成 value
//适合 m[]较大 w[]较小的情况
//返回布尔量
//传入货币种数 n,个数 m[],价值 w[]和目标值 value
#define MAXV 100000
int genfunc(int n,int* m,int* w,int value){
    int i,j,k,c;
    char r[MAXV];
    for (r[0]=i=1;i<=value;r[i++]=0);
    for (i=0; i< n; i++){
       for (j=0; j< w[i]; j++){}
           c=m[i]*r[k=j];
           while ((k+=w[i])<=value)</pre>
               if (r[k])
                   c=m[i];
               else if (c)
                   r[k]=1,c--;
           if (r[value])
               return 1;
       }
    }
    return 0;
}
```

4. 第 k 元素

```
//取第 k 个元素, k=0..n-1
//平均复杂度 0(n)
//注意 a[]中的顺序被改变
#define _{cp(a,b)} ((a)<(b))
typedef int elem_t;
elem_t kth_element(int n,elem_t* a,int k){
    elem_t t,key;
    int l=0,r=n-1,i,j;
   while (l<r){
        for (key=a[((i=l-1)+(j=r+1))>>1];i<j;){}
           for (j--;_cp(key,a[j]);j--);
           for (i++;_cp(a[i],key);i++);
           if (i<j) t=a[i],a[i]=a[j],a[j]=t;</pre>
        if (k>j) l=j+1;
       else r=j;
    }
    return a[k];
}
                                      5. 幻方构造
//幻方构造(l!=2)
#define MAXN 100
void dllb(int l,int si,int sj,int sn,int d[][MAXN]){
    int n,i=0,j=1/2;
    for (n=1; n<=l*l; n++) {
       d[i+si][j+sj]=n+sn;
       if (n%l){
           i=(i)?(i-1):(l-1);
           j=(j==l-1)?0:(j+1);
      }
       else
           i=(i==l-1)?0:(i+1);
   }
}
void magic_odd(int l,int d[][MAXN]){
   dllb(l,0,0,0,d);
}
void magic_4k(int l,int d[][MAXN]){
    int i,j;
    for (i=0;i<l;i++)
```

```
for (j=0;j<l;j++)
    d[i][j] = ((i\%4 = 0) | i\%4 = 3) \& (j\%4 = 0) | j\%4 = 3) | (i\%4 = 1) | i\%4 = 2) \& (j\%4 = 1) | j\%4 = 2))?(1*)
l-(i*l+j)):(i*l+j+1);
}
void magic_other(int l,int d[][MAXN]){
    int i, j, t;
   dllb(1/2,0,0,0,d);
   dllb(l/2,l/2,l/2,l*l/4,d);
    dllb(1/2,0,1/2,1*1/2,d);
   dllb(1/2,1/2,0,1*1/4*3,d);
    for (i=0;i<l/2;i++)
        for (j=0; j<1/4; j++)
            if (i!=1/4||j)
                t=d[i][j],d[i][j]=d[i+l/2][j],d[i+l/2][j]=t;
    t=d[1/4][1/4],d[1/4][1/4]=d[1/4+1/2][1/4],d[1/4+1/2][1/4]=t;
    for (i=0; i<1/2; i++)
        for (j=l-l/4+1;j<l;j++)
            t=d[i][j],d[i][j]=d[i+l/2][j],d[i+l/2][j]=t;
}
void generate(int l,int d[][MAXN]){
    if (1%2)
        magic_odd(l,d);
    else if (1\%4==0)
       magic_4k(l,d);
   else
       magic other(l,d);
}
                                   6. 模式匹配(kmp)
//模式匹配,kmp 算法,复杂度 0(m+n)
//返回匹配位置,-1表示匹配失败,传入匹配串和模式串和长度
//可更改元素类型,更换匹配函数
#define MAXN 10000
#define _{match(a,b)} ((a)==(b))
typedef char elem_t;
int pat_match(int ls,elem_t* str,int lp,elem_t* pat){
    int fail[MAXN]=\{-1\}, i=0, j;
    for (j=1;j<lp;j++){
        for (i=fail[j-1];i>=0&&!_match(pat[i+1],pat[j]);i=fail[i]);
        fail[j]=(_match(pat[i+1],pat[j])?i+1:-1);
    }
    for (i=j=0;i<ls&&j<lp;i++)
```

```
if (_match(str[i],pat[j]))
           j++;
       else if (j)
           j=fail[j-1]+1,i--;
    return j==lp?(i-lp):-1;
}
                                     7. 逆序对数
//序列逆序对数,复杂度 0(nlogn)
//传入序列长度和内容,返回逆序对数
//可更改元素类型和比较函数
#include <string.h>
#define MAXN 1000000
#define cp(a,b) ((a)<=(b))
typedef int elem_t;
elem_t _tmp[MAXN];
int inv(int n,elem_t* a){
   int l=n>>1, r=n-l,i,j;
   int ret=(r>1?(inv(l,a)+inv(r,a+l)):0);
   for (i=j=0;i<=l; tmp[i+j]=a[i],i++)
       for (ret+=j;j< r\&(i==l||!\_cp(a[i],a[l+j]));\_tmp[i+j]=a[l+j],j++);
   memcpy(a,_tmp,sizeof(elem_t)*n);
   return ret;
}
                                 8. 字符串最小表示
/*
   求字符串的最小表示
   输入: 字符串
   返回: 字符串最小表示的首字母位置(0...size-1)
*/
template <class T>
int MinString(vector <T> &str)
{
   int i, j, k;
   vector <T> ss(str.size() << 1);</pre>
   for (i = 0; i < str.size(); i ++) ss[i] = ss[i + str.size()] = str[i];
   for (i = k = 0, j = 1; k < str.size() && i < str.size() && j < str.size(); ) {
       for (k = 0; k < str.size() \&\& ss[i + k] == ss[j + k]; k ++);
       if (k < str.size()) {</pre>
           if (ss[i + k] > ss[j + k])
               i += k + 1;
           else j += k + 1;
           if (i == j) j ++;
       }
```

```
}
   return i < j? i : j;
}
                               9. 最长公共单调子序列
// 最长公共递增子序列, 时间复杂度 0(n^2 * logn), 空间 0(n^2)
/**
* n为a的大小, m为b的大小
* 结果在 ans 中
* "define _cp(a,b) ((a)<(b))"求解最长严格递增序列
*/
#define MAXN 1000
#define _{cp(a,b)} ((a)<(b))
typedef int elem_t;
elem_t DP[MAXN][MAXN];
int num[MAXN], p[1<<20];
int LIS(int n, elem_t *a, int m, elem_t *b, elem_t *ans){
   int i, j, l, r, k;
   DP[0][0] = 0;
   num[0] = (b[0] == a[0]);
   for(i = 1; i < m; i++) {
      num[i] = (b[i] == a[0]) || num[i-1];
      DP[i][0] = 0;
   for(i = 1; i < n; i++){
      if(b[0] == a[i] \&\& !num[0]) {
         num[0] = 1;
         DP[0][0] = i << 10;
      }
      for(j = 1; j < m; j++){
         for(k=((l=0)+(r=num[j-1]-1))>>1; l<=r; k=(l+r)>>1)
             if(_cp(a[DP[j-1][k]>>10], a[i]))
                l=k+1;
             else
                r=k-1;
         if(l < num[j-1] \&\& i == (DP[j-1][l] >> 10)){
             if(l >= num[j]) \ DP[j][num[j]++] = DP[j-1][l];
             else DP[j][l] = _cp(a[DP[j][l]>>10],a[i]) ? DP[j][l] : DP[j-1][l];
         }
```

 $if(b[j] == a[i]){$

```
for(k=((l=0)+(r=num[j]-1))>>1; l<=r; k=(l+r)>>1)
                if(_cp(a[DP[j][k]>>10], a[i]))
                   l=k+1;
                else
                   r=k-1;
             DP[j][l] = (i << 10) + j;
             num[j] += (l>=num[j]);
             p[DP[j][l]] = l ? DP[j][l-1] : -1;
         }
      }
   }
   for (k=DP[m-1][i=num[m-1]-1];i>=0;ans[i--]=a[k>>10],k=p[k]);
   return num[m-1];
}
                                   10. 最长子序列
//最长单调子序列,复杂度 0(nlogn)
//注意最小序列覆盖和最长序列的对应关系,例如
//"define cp(a,b) ((a)>(b))"求解最长严格递减序列,则
//"define _cp(a,b) (!((a)>(b)))"求解最小严格递减序列覆盖
//可更改元素类型和比较函数
#define MAXN 10000
#define _{cp(a,b)}((a)>(b))
typedef int elem t;
int subseq(int n,elem t* a){
   int b[MAXN],i,l,r,m,ret=0;
   for (i=0;i<n;b[l]=i++,ret+=(l>ret))
       for (m=((l=1)+(r=ret))>>1; l<=r; m=(l+r)>>1)
           if (_cp(a[b[m]],a[i]))
              l=m+1;
           else
               r=m-1;
   return ret;
}
int subseq(int n,elem_t* a,elem_t* ans){
   int b[MAXN],p[MAXN],i,l,r,m,ret=0;
   for (i=0; i< n; p[b[l]=i++]=b[l-1], ret+=(l>ret))
       for (m=((l=1)+(r=ret))>>1; l<=r; m=(l+r)>>1)
           if (_cp(a[b[m]],a[i]))
               l=m+1;
           else
               r=m-1;
```

```
for (m=b[i=ret];i;ans[--i]=a[m],m=p[m]);
    return ret;
}
                                  11. 最大子串匹配
//最大子串匹配,复杂度 0(mn)
//返回最大匹配值,传入两个串和串的长度,重载返回一个最大匹配
//注意做字符串匹配是串末的'\0'没有置!
//可更改元素类型,更换匹配函数和匹配价值函数
#include <string.h>
#define MAXN 100
#define max(a,b) ((a)>(b)?(a):(b))
#define _{match(a,b)} ((a)==(b))
#define value(a,b) 1
typedef char elem_t;
int str_match(int m,elem_t* a,int n,elem_t* b){
   int match[MAXN+1][MAXN+1],i,j;
   memset(match,0,sizeof(match));
   for (i=0; i< m; i++)
       for (j=0; j< n; j++)
           match[i+1][j+1] = max(max(match[i][j+1], match[i+1][j]),
                          (match[i][j]+_value(a[i],b[i]))*_match(a[i],b[j]));
   return match[m][n];
}
int str_match(int m,elem_t* a,int n,elem_t* b,elem_t* ret){
    int match[MAXN+1][MAXN+1],last[MAXN+1][MAXN+1],i,j,t;
   memset(match,0,sizeof(match));
   for (i=0; i< m; i++)
       for (j=0; j< n; j++){
           match[i+1][j+1]=(match[i][j+1]>match[i+1][j]?match[i][j+1]:match[i+1][j]);
           last[i+1][j+1]=(match[i][j+1]>match[i+1][j]?3:1);
           if ((t=(match[i][j]+_value(a[i],b[i]))*_match(a[i],b[j]))>match[i+1][j+1])
               match[i+1][j+1]=t,last[i+1][j+1]=2;
   for (;match[i][j];i-=(last[t=i][j]>1),j-=(last[t][j]<3))</pre>
       ret[match[i][j]-1]=(last[i][j]<3?a[i-1]:b[j-1]);
   return match[m][n];
}
                                   12. 最大子段和
//求最大子段和,复杂度 0(n)
```

```
typedef int elem_t;
elem t maxsum(int n,elem t* list){
    elem_t ret,sum=0;
    int i;
    for (ret=list[i=0];i<n;i++)</pre>
        sum=(sum>0?sum:0)+list[i],ret=(sum>ret?sum:ret);
    return ret:
}
elem_t maxsum(int n,elem_t* list,int& start,int& end){
    elem_t ret,sum=0;
    int s,i;
    for (ret=list[start=end=s=i=0];i<n;i++,s=(sum>0?s:i))
        if ((sum=(sum>0?sum:0)+list[i])>ret)
            ret=sum,start=s,end=i;
    return ret;
}
                                    13. 最大子阵和
//求最大子阵和,复杂度 0(n^3)
//传入阵的大小 m,n 和内容 mat[][]
//返回最大子阵和,重载返回子阵位置(maxsum=list[s1][s2]+...+list[e1][e2])
//可更改元素类型
#define MAXN 100
typedef int elem t;
elem_t maxsum(int m,int n,elem_t mat[][MAXN]){
    elem_t matsum[MAXN][MAXN+1], ret, sum;
    int i, j, k;
    for (i=0; i< m; i++)
        for (matsum[i][j=0]=0; j< n; j++)
           matsum[i][j+1]=matsum[i][j]+mat[i][j];
    for (ret=mat[0][j=0];j<n;j++)
        for (k=j;k< n;k++)
            for (sum=0, i=0; i< m; i++)
               sum=(sum>0?sum:0)+matsum[i][k+1]-matsum[i][j],ret=(sum>ret?sum:ret);
    return ret;
}
elem_t maxsum(int m,int n,elem_t mat[][MAXN],int& s1,int& s2,int& e1,int& e2){
    elem t matsum[MAXN][MAXN+1],ret,sum;
    int i,j,k,s;
    for (i=0;i< m;i++)
        for (matsum[i][j=0]=0; j< n; j++)
           matsum[i][j+1]=matsum[i][j]+mat[i][j];
```