

# Constructive Approximation Theory

## Assignment 9

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### 1 Aliasing of Trigonometric functions

$C_k = \cos(2k\pi x)$  For  $n \geq 1$  and  $0 \leq m \leq n$ ,

Let  $k = Nn \pm m$  where  $N = 1, 2, \dots$

For the  $n + 1$  grid, let  $x = \frac{j}{n}$  where  $j = 0, 1, \dots, n$

$$C_k = \cos(2(Nn \pm m)\pi \frac{j}{n}) = \cancel{\cos(2Nj\pi)} \cos(2m\pi \frac{j}{n}) \mp \cancel{\sin(2Nj\pi)} \sin(2m\pi \frac{j}{n}) = C_m$$

Thus,

$$C_m = C_{n \pm m} = C_{2n \pm m} = C_{3n \pm m} \dots$$

Similarly, for  $S_k = \sin(2k\pi x)$

$$S_k = \sin(2(Nn \pm m)\pi \frac{j}{n}) = \cancel{\sin(2Nj\pi)} \cos(2m\pi \frac{j}{n}) \pm \cancel{\cos(2Nj\pi)} \sin(2m\pi \frac{j}{n}) = \pm S_m$$

Thus

$$S_m = \pm S_{n \pm m} = \pm S_{2n \pm m} = \pm S_{3n \pm m} \dots$$

#### 1.1 Aliasing of Chebyshev polynomials

$T_k = \cos(k \arccos(x))$  For  $n \geq 1$  and  $0 \leq m \leq n$

Let  $k = 2Nn \pm m$  where  $N = 1, 2, \dots$

For the  $n + 1$  grid, let  $x = \cos(\frac{j\pi}{n})$  where  $j = 0, 1, \dots, n$

$$T_k = \cos(k \frac{j\pi}{n}) = \cos((2Nn \pm m) \frac{j\pi}{n}) = \cancel{\cos(2Nj\pi)} \cos(m\pi \frac{j}{n}) \mp \cancel{\sin(2Nj\pi)} \sin(m\pi \frac{j}{n}) = T_m$$

Thus,

$$T_m = T_{2n \pm m} = T_{4n \pm m} = T_{6n \pm m}$$

$$\text{For } k = 2n \pm m \implies m = \pm k \pmod{2n}$$

$$\text{Therefore, } T_k = T_m \text{ for } m = \pm k \pmod{2n}$$

We will use the following modular arithmetic identity for the next proof:

$$(a + b) \pmod{n} = (a \pmod{n} + b \pmod{n}) \pmod{n} \spadesuit$$

Now, for

$$\begin{aligned} m &= |(k + n - 1) \pmod{2n} - (n - 1)| \\ \implies \pm m &= (k + n - 1) \pmod{2n} - (n - 1) \\ \implies \pm m &= (k + n - 1) \pmod{2n} - (n - 1) \pmod{2n} \\ (\text{Because } n - 1 < 2n &\implies (n - 1) \pmod{2n} = n - 1) \\ \implies \pm m \pmod{2n} &= ((k + n - 1) \pmod{2n} - (n - 1) \pmod{2n}) \pmod{2n} \\ \implies \pm m &= ((k + n - 1) \pmod{2n} - (n - 1) \pmod{2n}) \pmod{2n} \\ (\text{Because } m < 2n &\implies m \pmod{2n} = m) \\ \implies \pm m &= (k + \cancel{n - 1} \cancel{- n + 1}) \pmod{2n} \\ (\text{Using the identity } \spadesuit) \\ \implies \pm m &= k \pmod{2n} \\ \implies m &= \pm k \pmod{2n} \end{aligned}$$

And for this relation, we have already proved that  $T_k = T_m$

## 2 Aliasing of Chebyshev coefficients

Let

$$f = \sum_{k=0}^{\infty} a_k T_k(x)$$

And, Let

$$c_0 = a_0 + a_{2n} + a_{4n} + a_{6n} + \dots$$

$$c_n = a_n + a_{3n} + a_{5n} + a_{7n} + \dots$$

For  $0 < k < n$

$$c_k = a_k + (a_{2n+k} + a_{4n+k} + a_{6n+k} + \dots) + (a_{2n-k} + a_{4n-k} + a_{6n-k} + \dots)$$

Let  $x_j$  be  $n + 1$  Chebyshev nodes. Multiplying  $T_k(x_j)$  with each  $c_k$  and considering that

$$T_m = T_{2n \pm m} = T_{4n \pm m} = T_{6n \pm m}$$

we get

$$\begin{aligned} c_0 T_0(x_j) &= (a_0 + a_{2n} + a_{4n} + a_{6n} + \dots) T_0(x_j) = a_0 T_0(x_j) + a_{2n} T_{2n}(x_j) + a_{4n} T_{4n}(x_j) \dots \\ c_n T_n(x_j) &= (a_n + a_{3n} + a_{5n} + a_{7n} + \dots) T_n(x_j) = a_n T_n(x_j) + a_{3n} T_{3n}(x_j) + a_{5n} T_{5n}(x_j) \dots \end{aligned}$$

$$\begin{aligned} c_k T_k(x_j) &= (a_k + (a_{2n+k} + a_{4n+k} + a_{6n+k} + \dots) + (a_{2n-k} + a_{4n-k} + a_{6n-k} + \dots)) T_k(x_j) \\ \implies c_k T_k(x_j) &= a_k T_k(x_j) + (a_{2n+k} T_{2n+k}(x_j) + a_{4n+k} T_{4n+k}(x_j) + \dots) + (a_{2n-k} T_{2n-k}(x_j) + a_{4n-k} T_{4n-k}(x_j) + \dots) \end{aligned}$$

Now considering the sum,

$$p_n(x_j) = \sum_{k=0}^n c_k T_k(x_j)$$

is a polynomial of degree  $n$ , and expanding the coefficients and rearranging the terms,

$$p_n(x_j) = \sum_{k=0}^{\infty} a_k T_k(x_j) = f(x_j)$$

Thus, at the Chebyshev nodes,  $x_j$

$$f(x_j) = p_n(x_j)$$

Thus,  $p_n$  is an  $n$ -degree interpolant of  $f$ , matching it at  $n + 1$  points, for the coefficients  $c_k$ . Aliasing of Chebyshev coefficients is hence proven.

### 3 Chebyshev approximation bounds

$$|a_k| \leq \frac{2M}{\rho^k}$$

Now,

$$\begin{aligned} |f - f_n| &= \left| \sum_{k=n+1}^{\infty} a_k T_k \right| \leq \left| \sum_{k=n+1}^{\infty} a_k \right| \leq \sum_{k=n+1}^{\infty} |a_k| \leq \sum_{k=n+1}^{\infty} \frac{2M}{\rho^k} = \frac{2M \rho^{-(n+1)}}{1 - \frac{1}{\rho}} = \frac{2M \rho^{-n}}{\rho - 1} \\ \implies \|f - f_n\| &\leq \frac{2M \rho^{-n}}{\rho - 1} \end{aligned}$$

Also,

$$\begin{aligned}
|f - p_n| &= \left| \sum_{k=0}^n (a_k - c_k) T_k + \sum_{k=n+1}^{\infty} a_k T_k \right| \leq \left| \sum_{k=0}^n (a_k - c_k) T_k \right| + \left| \sum_{k=n+1}^{\infty} a_k T_k \right| \\
\Rightarrow |f - p_n| &\leq \left| \sum_{k=0}^n (a_k - c_k) T_k \right| + \frac{2M\rho^{-n}}{\rho - 1} \leq \left| \sum_{k=0}^n (a_k - c_k) \right| + \frac{2M\rho^{-n}}{\rho - 1} \\
\Rightarrow |f - p_n| &\leq \left| \sum_{k=0}^n a_k - \sum_{k=0}^n c_k \right| + \frac{2M\rho^{-n}}{\rho - 1} = \left| \sum_{k=0}^n a_k - \sum_{k=0}^{\infty} a_k \right| + \frac{2M\rho^{-n}}{\rho - 1} \\
\Rightarrow |f - p_n| &\leq \left| - \sum_{k=n+1}^{\infty} a_k \right| + \frac{2M\rho^{-n}}{\rho - 1} \leq \frac{4M\rho^{-n}}{\rho - 1} \\
\Rightarrow \|f - p_n\| &\leq \frac{4M\rho^{-n}}{\rho - 1}
\end{aligned}$$