## Digital Logic and Circuit Paper Code: CS-102

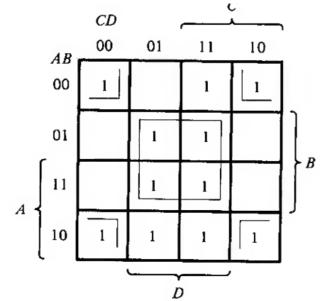
#### Outline

- ➤ Simplification of Boolean Functions using k-map
  - Prime and essential prime implicants
  - ➤ Sum of Minterm
  - ➤ Product of Maxterm

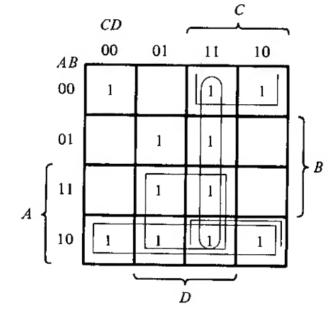
## Prime implicants and essential prime implicants

A prime implicant is a product term obtained by combining the maximum possible number of adjacent squares in the map.

If a minterm in a square is covered by only one prime implicant, that prime implicant is said to be essential.

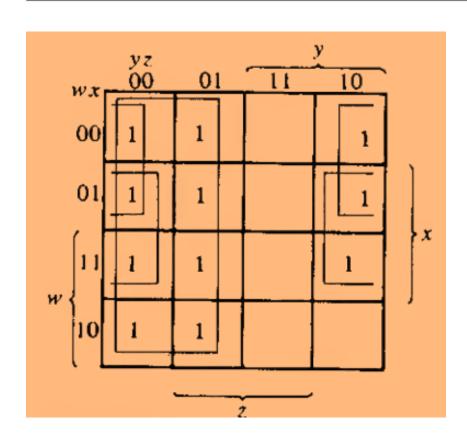


(a) Essential prime implicants BD and B'D'



(b) Prime implicants CD, B'C, AD, and AB'

## Simplify the Boolean function $F(w, x, y, z) = \sum (0, 1, 2, 4, 5, 6, 8, 9, 12, 13, 14)$



The simplified function is F= y' + w'z' + xz'

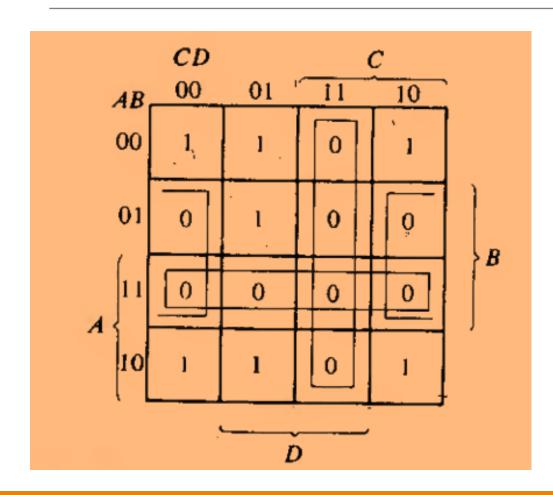
#### Simplify the Boolean function F = A'B'C' + B'CD' + A'BCD' + AB'C'

Number of adjacent Number of literals in a term in an n-variable map squares Z<sup>k</sup> n = 2n = 3n = 4n = 5n = 6n = 716 32 64

#### PRODUCT OF SUMS SIMPLIFICATION

- The procedure for obtaining a minimized function in product of sums follows from the basic properties of Boolean functions.
- The 1 s placed in the squares of the map represent the minterms of the function.
- The minterms not included in the function denote the complement of the function. From this we see that the complement of a function is represented in the map by the squares not marked by 1's.
- ➤ If we mark the empty squares by 0's and combine them into valid adjacent squares, we obtain a simplified expression of the complement of the function, i.e., of F '.

# Simplify the following Boolean function in (a) sum of products and (b) product of sum $F(A, B, C, D) = \sum (0, 1, 2, 5, 8, 9, 10)$

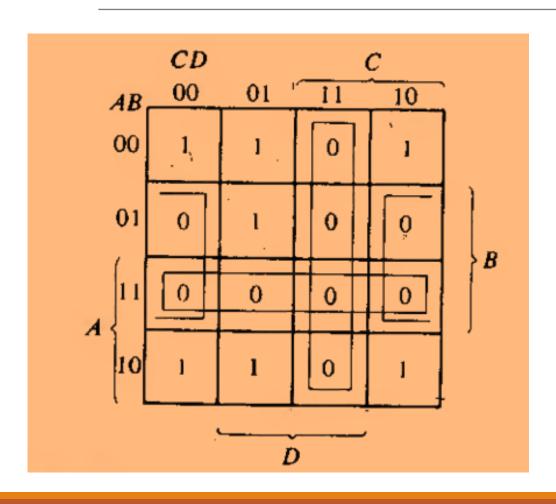


Combining the squares with I's gives the simplified function in sum of products:

$$F = B'D' + B'C' + A'C'D$$

If the squares marked with 0's are combined, as shown in the diagram, we obtain the simplified complemented function:

# Simplify the following Boolean function in (a) sum of products and (b) product of sum $F(A, B, C, D) = \sum (0, 1, 2, 5, 8, 9, 10)$



If the squares marked with 0's are combined, as shown in the diagram, we obtain the simplified complemented function:

Applying DeMorgan's theorem, we obtain the simplified function in product of sums:

$$F = (A' + B')(C' + D')(B' + D)$$

- The logical sum of the minterms associated with a Boolean function specifies the conditions under which the function is equal to 1.
- The function is equal to 0 for the rest of the minterms. This assumes that all the combinations of the values for the variables of the function are valid.
- In practice, there are some applications where the function is not specified for certain combinations of the variables.
- As an example, the four-bit binary code for the decimal digits has six combinations that are not used and consequently are considered as unspecified.
- In most applications, we simply don't care what value is assumed by the function for the unspecified minterms.
- For this reason, it is customary to call the unspecified minterms of a function don t-care conditions. These don't-care conditions can be used on a map to provide further simplification of the Boolean expression.

Real circuits don't always need to have an output defined for every possible input.

• For example, some calculator displays consist of 7-segment LEDs. These LEDs can display 2 <sup>7</sup> -1 patterns, but only ten of them are useful.

If a circuit is designed so that a particular set of inputs can never happen, we call this set of inputs a *don't care* condition.

They are very helpful to us in Kmap circuit simplification.

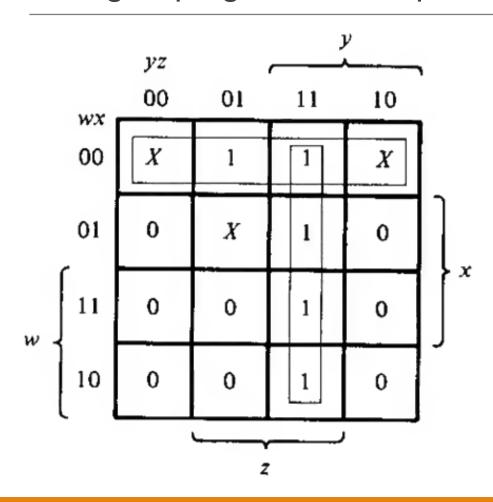
Example: Simplify the Boolean function  $F(w, x, y, z) = \sum (1, 3, 7, 11, 15)$  that has the don't-care conditions  $d(w, x, y, z) = \sum (0, 2, 5)$ 

In a Kmap, a don't care condition is identified by an X in the cell of the minterm(s) for the don't care inputs, as shown below.

In performing the simplification, we are free to include or ignore the X's when creating our groups.

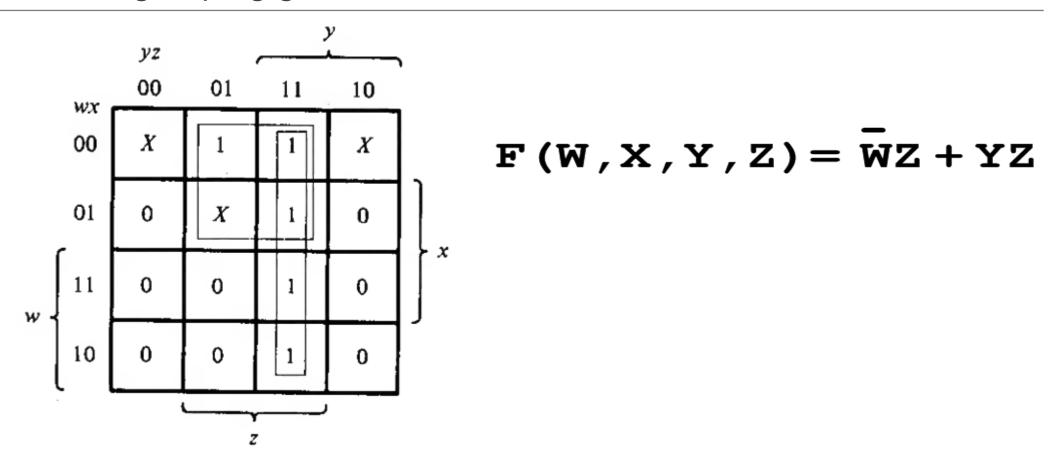
WX Y	z 00	01	11	10
00	×	1	1	×
01		×	1	
11			1	
10			1	
-			•	<b>1</b> 1 1

In one grouping in the Kmap below, we have the function:



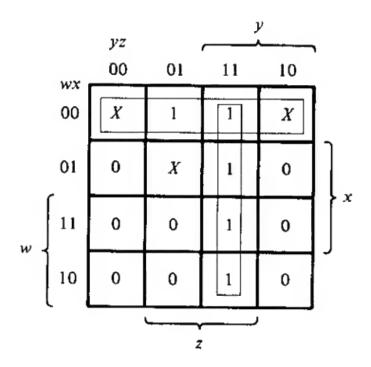
$$F(W,X,Y,Z) = \overline{WX} + YZ$$

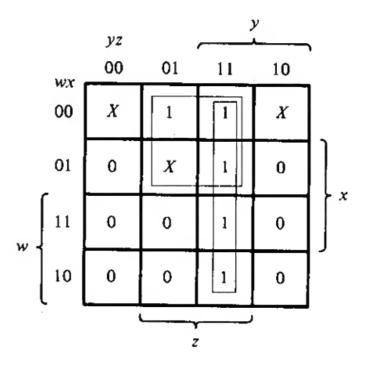
A different grouping gives us the function:



The truth table of: F(W,X,Y,Z) = WX+YZ

differs from the truth table of:  $\mathbf{F}(\mathbf{W}, \mathbf{X}, \mathbf{Y}, \mathbf{Z}) = \mathbf{WZ} + \mathbf{YZ}$ However, the values for which they differ, are the inputs for which we have don't care conditions.





### Suggested Reading

☐M. Morris Mano, Digital Logic and Computer Design, PHI.

### Thank you