Standard Code Library

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一切的开始

宏定义

数学

快速幂

1

• 注意乘法溢出

inline LLL fp(LLL a, LLL b, LLL Mod) {

```
LLL res = (Mod != 1);
        for (a %= Mod; b; b >>= 1, a = a * a % Mod)
3
            if (b & 1) res = res * a % Mod;
        return res;
5
    }
    GCD
    template <typename T>
    inline T gcd(T a, T b) {
        while (b){
           T t = b;
            b = a \% b;
            a = t;
        return a;
    }
    template <typename T>
11
    inline T lcm(T a, T b) { return a / gcd(a, b) * b; }
12
13
    template <typename T>
14
    inline T exgcd(T a, T b, T &x, T &y) {
        T m = 0, n = 1, t;
16
17
        x = 1, y = 0;
        while (b){
18
           t = m, m = x - a / b * m, x = t;
19
            t = n, n = y - a / b * n, y = t;
20
            t = b, b = a \% b, a = t;
21
22
        return a;
23
24
    }
    CRT
```

- 同余方程合并
- 返回最小正数解或最小非负解无解则返回-1

```
inline LL Crt(LL a1, LL a2, LL mod1, LL mod2) {
    LL u, v;
    LL g = exgcd(mod1, mod2, u, v);
```

```
if ((a2 - a1) % g)
5
            return -1;
        LL m12 = abs(lcm(mod1, mod2));
        LLL res = (((LLL)mod1 * ((LLL)u * ((a2 - a1) / g) % m12) % m12) + a1) % m12;
        return res <= 0 ? res + m12 : res; /* 求最小正数解还是非负解 */
   }
    线性筛
    struct primenumberlist{
    #define MAXN (10000000)
2
        int cnt, pri[10000000];
3
        bool np[MAXN + 10];
        primenumberlist(){
            np[1] = 1; cnt = 0;
            for (int i = 2; i <= MAXN; ++i) {</pre>
                if (!np[i]) pri[++cnt] = i;
                for (int j = 1; j <= cnt; ++j) {</pre>
                    LL t = pri[j] * i;
10
                    if (t > MAXN) break;
                    np[t] = 1;
12
13
                    if (!(i % pri[j])) break;
                }
14
            }
15
        }
   } prime;
17
    Φ欧拉函数
    template <typename T>
2
    inline T phi(T x) {
        T res = x;
3
        for (T i = 2; i * i <= x; ++i)</pre>
            if ((x % i) == 0) {
5
                res = res / i * (i - 1);
                while ((x \% i) == 0) x /= i;
        if (x > 1) res = res / x * (x - 1);
        return res;
10
   }
    Miller-Rabin 素性测试
       • n <= 10^{18}
    namespace MillerRabin {
1
        const LLL test[]={211,32511,937511,2817811,45077511,978050411,179526502211};
        inline bool isprime(LLL n) {
            if (n==13||n==19||n==73||n==193||n==407521||n==299210837ll)return 1;
5
            if (n <= 3) return n > 1;
            if (n <= 6) return n == 5;
            if (!(n & 1) || !(n % 3) || !(n % 5)) return 0;
            LLL d = n - 1; int t = 0;
10
            while (!(d & 1)) d >>= 1, ++t;
11
            for (LLL ai = 0, a = test[0]; ai < 7; ++ai, a = test[ai]) {</pre>
12
                if (a % n == 0) return 0;
13
                LLL v = fp(a, d, n); if (v == 1 | | v == n - 1) continue;
                LLL pre = v;
15
                for (int i = 1; i <= t; ++i) {</pre>
16
                    v = v * v % n;
17
                    if (v == 1)
18
                        if (pre != 1 && pre != (n - 1)) return 0; else break;
                    pre = v;
20
                if (v != 1) return 0;
22
23
24
            return 1;
```

```
25 }
26 }
```

Pollard-Rho 分解质因数

- 求 n 的一个非平凡因子
- 调用 $pollard_rho()$ 前先判断 n 的素性

```
namespace PollardRho{
        mt19937 mt(20011224);//19491001
        inline LLL pollard_rho(LLL n, LLL c) {
            LLL x = uniform_int_distribution < LL > (1, n - 1)(mt), y = x;
            LLL val = 1;
            for (int dep = 1;; dep <<= 1, x = y, val = 1) {</pre>
                 for (int stp = 1; stp <= dep; ++stp) {</pre>
                     y = (y * y + c) % n;
                     val = val * abs(x - y) % n;
10
                     if ((stp & 127) == 0) {
11
12
                         LLL d = gcd(val, n);
                         if (d > 1) return d;
13
                     }
                 }
15
16
                 LLL d = gcd(val, n);
                 if (d > 1) return d;
17
            }
18
        }
19
20
21
        //接口根据题意重写
        vector<LLL> factor;
22
        void getfactor(LLL x, LLL c = 19260817) {
23
            if (MillerRabin::isprime(x)) {factor.emplace_back(x); return;}
24
            LLL p = x;
25
            while (p == x) p = pollard_rho(x, c--);
26
27
            getfactor(p); getfactor(x / p);
28
        inline LLL ask(LLL x) {
29
            factor.clear();
30
31
            while (!(x & 1)) x >>= 1, factor.emplace_back(2);
            if (x > 1) getfactor(x);
32
            return factor.size();
        }
34
   }
35
```

组合数

- 数较小模数为较大质数求逆元
- - 如果模数固定可以 O(n) 预处理阶乘的逆元
- 数较大模数为较小质数用 Lucas 定理
- –

$$C_n^m \equiv C_{\lfloor \frac{n}{p} \rfloor}^{\lfloor \frac{m}{p} \rfloor} * C_{n \bmod p}^{m \bmod p} (mod \ p)$$

• 数较大模数较小用 exLucas 定理求 $C_n^m mod P$

exLucas

- O(P log P)
- 不要求 P 为质数

```
namespace EXLUCAS {
inline LL idxp(LL n, LL p) {
LL nn = n;
while (n > 0) nn -= (n % p), n /= p;
```

```
return nn / (p - 1);
        }
6
        LL facp(LL n, LL p, LL pk) {
8
            if (n == 0) return 1;
            LL res = 1;
10
            if (n >= pk) {
11
                 LL t = n / pk, k = 1, els = n - t * pk;
12
                 for (LL i = 1; i <= els; ++i) if (i % p) k = k * i % pk;
13
                 res = k;
14
                 for (LL i = els + 1; i < pk; ++i) if (i % p) k = k * i % pk;</pre>
15
16
                 res = res * fp(k, n / pk, pk) % pk;
17
            }
            else for (LL i = 1; i <= n; ++i) if (i % p) res = res * i % pk;
18
19
            return res * facp(n / p, p, pk) % pk;
20
21
        inline LL exlucas(LL n, LL m, LL p, LL pk, LL k) {
22
23
            LL a = facp(n, p, pk) * fp(facp(n - m, p, pk) * facp(m, p, pk) % pk, pk / p * (p - 1) - 1, pk) % pk;
            LL b = idxp(n, p) - idxp(m, p) - idxp(n - m, p);
24
            if (b >= k) return 0;
25
            while (b--) a *= p;
26
            return a % pk;
27
        }
29
        /* 接口 */ inline LL exlucas(LL n, LL m, LL p) {
30
31
            LL a = 0, b = 1;
            for (LL i = 2; i * i <= p; ++i) {
32
                 if (p % i) continue;
                 LL t = 0, pk = 1;
34
                 while (p \% i == 0) ++t, p /= i, pk *= i;
35
                 a = Crt(a, exlucas(n, m, i, pk, t), b, pk);
36
                 b \star = pk;
37
            }
            return (p > 1) ? Crt(a, exlucas(n, m, p, p, 1), b, p) : a;
39
40
    }
41
```

二维计算几何

- Point 直接支持整型和浮点型
- 部分函数可以对整型改写
- 多边形 (凸包) 按逆时针存在下标 1..n

点向量基本运算

```
1
    template <typename T>
    struct Point {
2
        T x, y;
3
        Point() {}
        Point(T u, T v) : x(u), y(v) {}
5
        Point operator+(const Point &a) const { return Point(x + a.x, y + a.y); }
        Point operator-(const Point &a) const { return Point(x - a.x, y - a.y); }
        Point operator*(const T &a) const { return Point(x * a, y * a); }
        T operator*(const Point &a) const { return x * a.x + y * a.y; }
        T operator%(const Point &a) const { return x * a.y - y * a.x; }
10
11
        double len() const { return hypot(x, y); }
        double operator^(const Point &a) const { return (a - (*this)).len(); }
12
        double angle() const { return atan2(y, x); }
14
        bool id() const { return y < 0 || (y == 0 && x < 0); }
        bool operator<(const Point &a) const { return id() == a.id() ? (*this) % a > 0 : id() < a.id(); }</pre>
15
16
    typedef Point<double> point;
17
   #define sqr(x) ((x) * (x))
19
    const point O(0, 0);
20
    const double PI(acos(-1.0)), EPS(1e-8);
21
    inline bool dcmp(const double &x, const double &y) { return fabs(x - y) < EPS; }
```

```
inline int sgn(const double &x) \{ return fabs(x) < EPS ? 0 : ((x < 0) ? -1 : 1); \}
   inline double mul(point p1, point p2, point p0) { return (p1 - p0) % (p2 - p0); }
   位置关系
   inline bool in_same_seg(point p, point a, point b) {
       if (fabs(mul(p, a, b)) < EPS) {</pre>
2
           if (a.x > b.x) swap(a, b);
           return (a.x <= p.x && p.x <= b.x && ((a.y <= p.y && p.y <= b.y) || (a.y >= p.y && p.y >= b.y)));
       } else return 0;
5
   }
   inline bool is_right(point st, point ed, point a) {
       return ((ed - st) % (a - st)) < 0;
10
11
   inline point intersection(point s1, point t1, point s2, point t2) {
12
13
       return s1 + (t1 - s1) * (((s1 - s2) % (t2 - s2)) / ((t2 - s2) % (t1 - s1)));
14
15
   inline bool parallel(point a, point b, point c, point d) {
16
17
       return dcmp((b - a) % (d - c), 0);
18
19
   inline double point2line(point p, point s, point t) {
20
       return fabs(mul(p, s, t) / (t - s).len());
21
22
23
   inline double point2seg(point p, point s, point t) {
24
       25
26
   多边形
   求多边形面积
   inline double area(int n, point s[]) {
2
       double res = 0;
       s[n + 1] = s[1];
       for (int i = 1; i <= n; ++i)</pre>
          res += s[i] % s[i + 1];
       return fabs(res / 2);
   }
   判断点是否在多边形内
       • 特判边上的点
       ● 使用了 a[1]...a[n+1] 的数组
   inline bool in_the_area(point p, int n, point area[]) {
       bool ans = 0; double x;
2
       area[n + 1] = area[1];
       for (int i = 1; i <= n; ++i) {</pre>
           point p1 = area[i], p2 = area[i + 1];
           if (in_same_seg(p, p1, p2)) return 1; //特判边上的点
           if (p1.y == p2.y) continue;
           if (p.y < min(p1.y, p2.y)) continue;</pre>
           if (p.y >= max(p1.y, p2.y)) continue;
           ans ^{=} (((p.y - p1.y) * (p2.x - p1.x) / (p2.y - p1.y) + p1.x) > p.x);
11
       return ans;
12
   }
13
   凸包
       ● Andrew 算法
```

O(n log n)

● 注意是否应该统计凸包边上的点

```
inline bool pcmp1(const point &a, const point &b) { return a.x == b.x ? a.y < b.y : a.x < b.x; }</pre>
1
    inline int Andrew(int n, point p[], point ans[]) { //ans[] 逆时针存凸包
3
        sort(p + 1, p + 1 + n, pcmp1);
        int m = 0;
        for (int i = 1; i <= n; ++i) {</pre>
            while (m > 1 && mul(ans[m - 1], ans[m], p[i]) < 0) --m; //特判凸包边上的点
           ans[++m] = p[i];
        int k = m;
10
        for (int i = n - 1; i >= 1; --i) {
11
            while (m > k && mul(ans[m - 1], ans[m], p[i]) < 0) --m; //特判凸包边上的点
12
           ans[++m] = p[i];
13
15
        return m - (n > 1); //返回凸包有多少个点
   }
16
    凸包直径·平面最远点对
       • 旋转卡壳算法

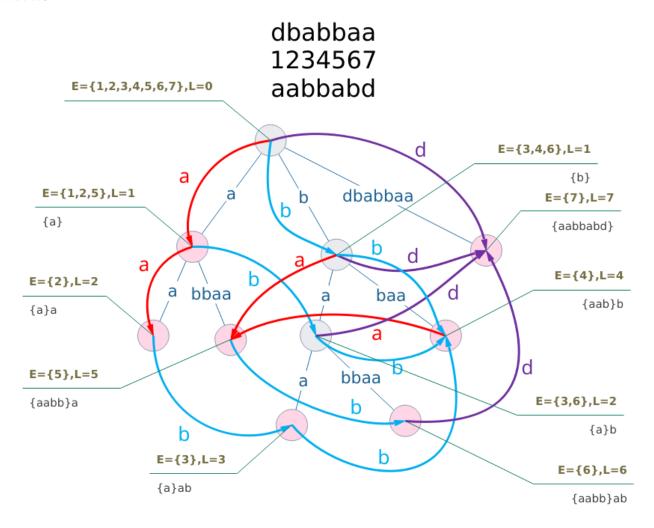
    O(n)

       • 凸包的边上只能有端点, 否则不满足严格单峰
       ● 使用了 a[1]...a[n+1] 的数组
    inline double Rotating_Caliper(int n, point a[]) {
        a[n + 1] = a[1];
        double ans = 0;
3
        int j = 2;
4
        for (int i = 1; i <= n; ++i) {</pre>
5
            while (fabs(mul(a[i], a[i + 1], a[j])) < fabs(mul(a[i], a[i + 1], a[j + 1]))) j = (j % n + 1);</pre>
            ans = \max(ans, \max((a[j] ^ a[i]), (a[j] ^ a[i + 1])));
        return ans;
   }
10
    平面最近点对
       分治+归并
       • O(n log n)
    namespace find_the_closest_pair_of_points {
        const int N = 200010; //maxn
2
        inline bool cmp1(const point &a, const point &b) { return a.x < b.x \mid \mid (a.x == b.x \&\& a.y < b.y); }
3
        inline bool operator (const point &a, const point &b) { return a.y > b.y | | (a.y == b.y \&\& a.x > b.x); }
5
        point a[N], b[N];
        double ans;
7
        inline void upd(const point &i, const point &j) { ans = min(ans, i ^ j); }
8
        void find(int l, int r) {
10
            if (l == r) return;
            if (l + 1 == r) {
12
13
                if (a[l] > a[r]) swap(a[l], a[r]);
                upd(a[l], a[r]); return;
14
15
           int mid = (l + r) >> 1;
           double mx = (a[mid + 1].x + a[mid].x) / 2;
17
            find(l, mid); find(mid + 1, r);
18
            int i = l, j = mid + 1;
19
            for (int k = 1; k \le r; ++k) b[k] = a[((j > r) | | (i \le mid \&\& a[j] > a[i])) ? (i++) : (j++)];
20
            for (int k = l; k <= r; ++k) a[k] = b[k];</pre>
21
            int tot = 0;
22
            for (int k = l; k \le r; ++k) if (fabs(a[k].x - mx) \le ans) {
23
                for (int j = tot; j >= 1 && (a[k].y - b[j].y <= ans); --j) upd(a[k], b[j]);
24
25
                b[++tot] = a[k];
            }
26
27
        }
```

```
//接口
29
30
        inline double solve(int n, point ipt[]){
             ans = 0x3f3f3f3f3f3f3f3f3f1l; //max distance
31
             for (int i = 1; i <= n; ++i) a[i] = ipt[i];</pre>
32
             sort(a + 1, a + 1 + n, cmp1);
33
             find(1, n);
34
35
             return ans;
        }
36
    }
    员
    三点垂心
    inline point geto(point p1, point p2, point p3) {
        double a = p2.x - p1.x;
2
        double b = p2.y - p1.y;
double c = p3.x - p2.x;
3
        double d = p3.y - p2.y;
        double e = sqr(p2.x) + sqr(p2.y) - sqr(p1.x) - sqr(p1.y);
        double f = sqr(p3.x) + sqr(p3.y) - sqr(p2.x) - sqr(p2.y);
return {(f * b - e * d) / (c * b - a * d) / 2, (a * f - e * c) / (a * d - b * c) / 2};
    }
    最小覆盖圆
        ● 随机增量 O(n)
    inline void min_circlefill(point &o, double &r, int n, point a[]) {
1
        mt19937 myrand(20011224); shuffle(a + 1, a + 1 + n, myrand); //越随机越难 hack
2
        o = a[1];
3
        r = 0;
4
        for (int i = 1; i <= n; ++i) if ((a[i] ^ o) > r + EPS) {
5
             o = a[i];
             r = 0;
             for (int j = 1; j < i; ++j) if ((o ^ a[j]) > r + EPS) {
                 o = (a[i] + a[j]) * 0.5;
                 r = (a[i] ^ a[j]) * 0.5;
10
                 for (int k = 1; k < j; ++k) if ((o ^ a[k]) > r + EPS) {
11
                     o = geto(a[i], a[j], a[k]);
12
                      r = (o ^ a[i]);
13
                 }
14
15
             }
        }
16
   }
17
```

字符串

后缀自动机



杂项

STL

copy

```
template <class InputIterator, class OutputIterator>
```

OutputIterator copy (InputIterator first, InputIterator last, OutputIterator result);