DATA STRUCTURE AND ALGORITHM

CLASS 3

Seongjin Lee

Updated: 2017-03-06 DSA_2017_03

insight@gnu.ac.kr http://resourceful.github.io Systems Research Lab. GNU



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ARRAY

an **array** is a set of pairs, **<index**, **value>**, such that each index that is defined has a value associated with it

- "a consecutive set of memory locations" in C
- $\, igcup \,$ logical order is the same as physical order

operations

- creating a new array
- retrieves a value
- stores a value
- insert a value into array delete a value at the array

A one-dimensional array in C is declared implicitly by appending brackets to the name of a variable

```
int list[5], *plist[5];
```

arrays start at index o in C

consider the implementation of one-dimensional arrays

```
int list[5];
```

- allocates five consecutive memory locations
- each memory location is large enough to hold a single integer
- base address address of the first element

```
list = &list[0]
```

Variable	Memory Address	
&list[0]	base address = α	
&list[1]	α + sizeof(int)	
&list[2]	$\alpha + 2 \cdot \text{sizeof(int)}$	
&list[3]	$\alpha + 3 \cdot \text{sizeof(int)}$	
&list[4]	$\alpha + 4 \cdot \text{sizeof(int)}$	

&list[i] in a C programs

 $\, \bigcirc \,$ C interprets it as a pointer to an integer or its value

```
int *list1; // pointer variable to an int
```

```
(list2+i) equals &list2[i], and *(list2+i) equals list2[i]
```

○ regardless of the type of the array list2

consider the way C treats an array when it is a parameter to a function

- the parameter passing is done using call-by-value in C
- but array parameters have their values altered

Array: Example 2.1

```
#define MAX_SIZE 100
float sum(float [], int);
float input[MAX_SIZE], answer;
int i;
void main(void) {
   for(i = 0; i < MAX_SIZE; i++) input[i] = i;
       answer = sum(input, MAX_SIZE);
   printf{"The sum is: %f\n", answer);
float sum(float list[], inst n) {
   int i;
   float tempsum = 0;
   for(i = 0; i < n; i++)
       tempsum += list[i];
   return tempsum;
```

Pointer Variable stores address

- & : Starting address of allocated variable
- *: Value stored on the address of the pointer variable

Do Not!

 pointer variable is not referencing an address, so cannot store a value

```
int *ptr;
* ptr = 100;
```

the data type must equal

```
double Pi = 3.14;
int *pPi = Π
```

cannot dereference a non-pointer variable

```
int num;
*num = 100;
```

Do!

it is recommended to initialize a pointer value with NULL ('\o')
 See the example

```
#include <stdlib.h>
```

- void *malloc(size_t size); // allocates size bytes of memory and returns
 a pointer to the allocated memory
- void *free(void *ptr); // frees allocation that were created via the preceding allocation function
- void *calloc(size_t count, size_t size); // contiguously allocates
 enough sapce for count objects that are size bytes of memory each
 and returns a pointer to the allocated memory. The allocated memory
 is filled with bytes of value zero.
- void *realloc(void *ptr, size_t size); // change the size of the
 allocation pointed to by ptr to size, and returns ptr

Array: Ex 2.2, 1-dimensional array addressing

```
int one[] = \{0, 1, 2, 3, 4\};
```

write a function that prints out both the address of the i^{th} element of the array and the value found at this address

Array: Ex 2.2 One dimensional array accessed by address

```
// codes/array_address.c
    #include <stdio h>
    void print1(int *ptr,int rows) {
        int i;
        printf("Address\t\tContents\n"):
       for(i=0;i<rows;i++)</pre>
            printf("%8u\t%5d\n", (unsigned int)ptr+i, *(ptr+i));
        printf("\n");
10
11
    int main() {
12
      int one[] = \{0, 1, 2, 3, 4\};
      print1(one, 5);
14
      return 0;
15
16
```

One-dimensional array accessed by address

- o address of i th element ptr + i
- obtain the value of the i th value *(ptr + i)

Address Contents 1518325632 0 1518325633 1 1518325634 2 1518325635 3 1518325636 4

One-dimensional array addressing

- the addresses increase by two on an Intel 386 machine
- O Example shown is the result of Mac OS X on Intel Core i5 Machine

STRUCTURES AND UNIONS

struct

- structure or record
- the mechanism of grouping data
- permits the data to vary in type

collection of data items where

o each item is identified as to its type and name

```
struct {
   char name[10];
   int age;
   float salary;
} person;
```

creating a variable

- whose name is person and
- has three fields
 - 1. a name that is a character array
 - 2. an integer value representing the age of the person
 - 3. a float value representing the salary of the individual

use of the . as the structure member operator

```
strcpy(person.name, "james");
person.age = 30;
person.salaray = 35000;
```

select a particular member of the structure

typedef statement

create our own structure data type

type 1

```
typedef struct human_being {
   char name[10];
   int age;
   float salary;
};
```

type 2

```
typedef struct {
    char name[10];
    int age;
    float salary;
} human_being;
```

human_being

 $\, \bigcirc \,$ the name of the type defined by the structure definition

```
human_being person1, person2;

if(strcmp(person1.name, person2.name))
    printf("The two people do not have the same name\n");
else
    printf("The two people have the same name\n");
```

Structures and Unions: Assignment

assignment

permits structure assignment in ANSI C

```
person1 = person2;
```

 but, in most earlier versions of C assignment of structures is not permitted

```
strcpy(person1.name,person2.name);
person1.age=person2.age;
person1.salary=person2.salary;
```

Structures and Unions: Equality or Inequality

equality or inequality

- cannot be directly checked
- Example function to check equality of struct

```
int humans_equal(human_being person1, human_being person2) {
    if(strcmp(person1.name,person2.name))
        return FALSE;
    if(person1.age!=person2.age)
        return FALSE;
    if(person1.salary!=person2.salary)
        return TALSE;
    return TRUE;
}
```

Structures and Unions: Embedding Structure

Embedding of a structure within a structure

```
typedef struct {
   int month;
   int day;
   int year;
} date;

typedef struct human_being {
   char name[10];
   int age;
   float salary;
   date dob; // embedded structure
};
```

Structures and Unions: Embedding Structure

Ex. A person boar on Feb 14 1992

```
person1.dob.month = 2;
person1.dob.day = 14;
person1.dob.year = 1992;
```

Structures and Unions: Unions

Unions

- similar to a structure, but
- the fields of a union must share their memory space
- only one field of the union is "active" at any given time

Structures and Unions: Unions

```
typedef struct sex_type {
   enum tag_field {female, male} sex;
   union {
       int children;
       int beard; } u;
};
typedef struct human_being {
   char name[10];
   int age;
   float salary;
   date dob;
   sex_type sex_info;
};
human_being person1,person2;
```

Structures and Unions: Unions

Assign values to person1 and person2

```
persone1.sex_info.sex = male;
person1.sex_info.u.beard = FALSE; /* FALSE: 0 */
```

and

```
person2.sex_info.sex = female;
person2.sex_info.u.children = 4;
```

Structures and Unions: Internal Implementation of Structure

```
struct {
   int i, j; float a, b;
```

or

```
struct {
   int i; int j; float a; float b;
};
```

stored in the same way

increasing address locations in the order specified in the structure definition

size of an object of a struct or union type

• the amount of storage necessary to represent the largest component

Data Structure and Algorithm

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Structures and Unions: Self-referential structures

- one or more of its components is a pointer to itself
- usually require dynamic storage management routines to explicitly obtain and release memory

```
typedef struct list {
   char data;
   list *link;
};
```

the value of link

o address in memory of an instance of list or null pointer

Structures and Unions: Self-referential structures

```
list item1, item2, item3;
item1.data = 'a';
item2.data = 'b';
item3.data = 'c';
item1.link = item2.link = item3.link = NULL;
```

Structures and Unions: Self-referential structures

```
list item1, item2, item3;
item1.data = 'a';
item2.data = 'b';
item3.data = 'c';
item1.link = item2.link = item3.link = NULL;
```

attach these structures together

```
item1.link = &item2;
item2.link = &item3;
```



The Sparse Matrix: representing matrix by using array

- \bigcirc *m* by *n* (*m* rows, *n* columns)
- use two-dimensional array
- space complexity S(m, n) = m * n

	col0	col1	col2
row0	[-27]	3	4]
row1	6	82	-2
row2	109	-64	11
row3	12	8	9
row4	48	27	47]

A[6,6]

	col0	col1	col2	col3	col4	col5
row0 row1 row2 row3 row4	г 15	0	0	22	0	15 ๅ
row1	0	11	3	0	0	0
row2	0	0	0	-6	0	0
row3	0	0	0	0	0	0
row4	91	0	0	0	0	0
row5		0	28	0	0	0]

common characteristics

- most elements are zero's
- inefficient memory utilization

solutions

- store only nonzero elements
- o using the triple <row, col, value>
- must know
 - the number of rows
 - the number of columns
 - the number of non-zero elements

```
#define MAX_TERMS 101

typedef struct {
   int col;
   int row;
   int value;
} term;
```

- a[o].row: the number of rows
- a[o].col: the number of columns
- a[o].value: the total number of non-zoros
- choose row-major order

The Sparse Matrix: As Tripels

	row	col	value
a[o]	6	6	8
a[1]	O	O	15
a[2]	O	3	22
a[3]	0	5	- 15
a[4]	1	1	11
a[5]	1	2	3
a[6]	2	3	-6
a[7]	4	O	91
a[8]	5	2	28

- space complexity (*variable space requirement)
 S(m,n) = 3 * t where
 t: the number of non-zero's
- independent of the size of rows and columns

The Sparse Matrix: Transpose

Transpose a matrix

- interchange rows and columns
- o move a[i][j] to a[j][i]

	row	col	value
b[o]	6	6	8
b[1]	O	O	15
b[2]	O	4	91
b[3]	1	1	11
b[4]	2	1	3
b[5]	2	5	28
b[6]	3	O	22
b[7]	3	2	-6
b[8]	5	O	- 15

```
algorithm BAD_TRANS
for each row i
   take element <i,j,value>;
   store it as element <j,i,value> of the transpose;
end;
```

The problem

```
algorithm BAD_TRANS
for each row i
   take element <i,j,value>;
   store it as element <j,i,value> of the transpose;
end;
```

The problem

data movement

```
algorithm BAD_TRANS
for each row i
   take element <i,j,value>;
   store it as element <j,i,value> of the transpose;
end;
```

The problem

data movement

```
algorithm TRANS
for all elements in column j
   place element <i,j,value> in element <j,i,value> end;
```

The problem

```
algorithm BAD_TRANS
for each row i
   take element <i,j,value>;
   store it as element <j,i,value> of the transpose;
end;
```

The problem

data movement

```
algorithm TRANS
for all elements in column j
   place element <i,j,value> in element <j,i,value> end;
```

The problem

unnecessary loop for each column

The Sparse Matrix: Transpose

```
void transpose(term a[], term b[]) {
   int n, i, j, currentb;
   n = a[0].value;
   b[0].row = a[0].col;
   b[0].col = a[0].row;
   b[0].value = n;
   if(n > 0){}
       currentb = 1;
       for (i = 0; i < a[0].col; i++)
           for(i = 1: i \le n: i++)
              if(a[i].col == i) {
                b[currentb].row = a[j].col;
                b[currentb].col = a[j].row;
                b[currentb].value = a[j].valuel;
                cuurentb++;
```

Space and Time complexity?

The Sparse Matrix: Transpose

Complexity

- space: 3× t
- \bigcirc time: $O(cols \times t)$

create better algorithm by using a little more storage

- row_terms the number of element in each row
- starting_pos the starting point of each row

The Sparse Matrix: Better Transpose I

```
void fast_transpose(term a[], term b[]){
   int row_terms[MAX_COL];
   int starting_pos[MAX_COL];
   int i, j;
   int num_cols = b[0].row = a[0].col;
   int num_terms = b[0].value = a[0].value;
   b[0].col = a[0].row:
   if(num_terms > 0) {
       for(i = 0; i < num\_cols; i++)
           row_terms[i] = 0;
       for(i = 1; i <= num_terms; i++)</pre>
           row_terms[a[i].col]++;
       starting_pos[0] = 1;
```

The Sparse Matrix: Better Transpose II

The Sparse Matrix: Better Transpose

```
[0] [1] [2] [3] [4] [5] row_terms= 2 1 2 2 0 1 starting_pos= 1 3 4 6 8 8 complexity
```

complexity

- space: 3× t + extra
- \bigcirc time: O(cols + t)

REPRESENTATION OF MULTIDIMEN-

SIONAL ARRAYS

internal representation of multidimensional arrays

- how to state n-dimensional array into 1-dimensional array?
- how to retrieve arbitrary element in an array a[upper⁰][upper¹] \cdots [upperⁿ⁻¹]

The number of elements in the array $\prod_{i=0}^{n-1} upper^i$ Ex. a[10][10]

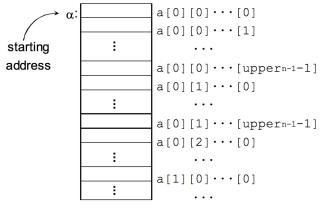
 $0 10 \times 10 \times 10 = 1000 \text{ (units)}$

represent multidimensional array by

in what order?

row-major-order

store multidimensional array by rows



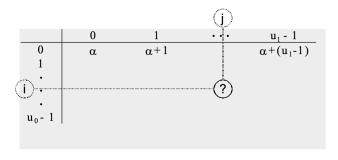
how to retreive values

- starting-address + offset-value
- assume α: starting-address

1-dimensional array a[u⁰]

```
&a[0] : \alpha &a[1] : \alpha + 1 : \alpha + 1 : \alpha &a[u^0-1] : \alpha + (u^0-1) &a[i] : \alpha + i
```

2-dimensional array a[u⁰][u¹]



&a[i][j] =
$$\alpha + i \cdot u^1 + j$$

2-dimensional array a[u⁰][u¹][u²]

&a[i][j][k] =
$$\alpha + i \cdot u^1 \cdot u^2 + j \cdot u^2 + k$$

= $\alpha + u^2[i \cdot u^1 + j] + k$

general case array $a[u^0][u^1]$ $cdots[u^{n-1}]$

&
$$a[i_0][i_1][i_{n-1}]$$

$$= \alpha + \sum_{j=0}^{n-1} i_j \cdot a_j \begin{cases} a_j = \prod_{k=j+1}^{n-1} u_k & \text{for } j < n-1 \\ a_{n-1} = 1 & \text{for } j = n-1 \end{cases}$$