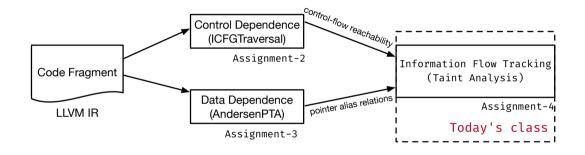
## **Information Flow Tracking**

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### **Today's Class**



### What is Taint Analysis?

- Taint analysis aims to reason about the control and data dependence from a source (statement/node) to a sink (statement/node).
- Taint analysis can also be seen as information flow tracking analysis.
  - Static taint analysis: taint tracking at compile time (this subject)
  - dynamic taint analysis: taint tracking during runtime.

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#### Why learn Taint Analysis?

- Detect information leakage
  - A sensitive data stored in a heap object and manipulated by pointers can be passed around and stored to an unchecked memory (untrusted third-party APIs)
- Detect code vulnerability
  - There is a vulnerability if an unchecked tainted source (e.g., return value from an untrusted third party function) flows into one of the following sinks, where the tainted variable being used as
    - a parameter passing to a sensitive function or
    - a bound access (array index) or
    - a termination condition (loop condition)
      - . . .

## **How to Perform Static Taint Analysis?**

Let us use what we have learned about control- and data-dependence to develop an information flow checker to validate tainted flows from a source to a sink.

- A source v<sub>src</sub>@s<sub>src</sub> is a tuple consisting of a variable v<sub>src</sub> and a statement s<sub>src</sub> where v<sub>src</sub> is defined.
- A sink v<sub>snk</sub>@s<sub>snk</sub> is also a tuple consisting of a variable v<sub>snk</sub> and a statement s<sub>snk</sub> where v<sub>snk</sub> is used.
- In SVF, variables  $v_{src}$  and  $v_{snk}$  are PAGNodes. Statements  $s_{src}$  and  $s_{snk}$  are ICFGNodes.

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- A sink v<sub>snk</sub>@s<sub>snk</sub> is also a tuple consisting of a variable v<sub>snk</sub> and a statement s<sub>snk</sub> where v<sub>snk</sub> is used.
- In SVF, variables v<sub>src</sub> and v<sub>snk</sub> are PAGNodes. Statements s<sub>src</sub> and s<sub>snk</sub> are ICFGNodes.
- Given a tainted source v<sub>src</sub>@s<sub>src</sub>, we say that a sink v<sub>snk</sub>@s<sub>snk</sub> is also tainted
  if both of the following two conditions satisfy:
  - (1) s<sub>src</sub> reaches s<sub>snk</sub> on the ICFG (Assignment 2), and
  - (2)  $v_{src}$  is aliased with  $v_{snk}$ , i.e.,  $pts(src) \cap pts(snk) \neq \emptyset$  (Assignment 3)

#### **Example 1**

```
int main(){
char* secretToken = tgetstr();  // source
char* a = secretToken;
char* b = a;
broadcast(b);  // sink
}
```

What is the tainted flow?

#### Example 1

```
int main(){
char* secretToken = tgetstr();  // source
char* a = secretToken;
char* b = a;
broadcast(b);  // sink
}
```

#### What is the tainted flow?

- Line 2 reaches Line 5 along the ICFG (control-dependence holds)
   secretToken and b are aliases (data-dependence holds)
- Both control-dependence and data-dependence hold. Therefore, secretToken@Line 2 flows to b@Line 5.

#### Example 2

```
int main(){
char* secretToken = tgetstr(...); // source
char* a = secretToken;
char* b = a;
char* publicToken = "hello";
broadcast(publicToken); // sink
}
```

#### Example 2

```
int main(){
char* secretToken = tgetstr(...); // source
char* a = secretToken;
char* b = a;
char* publicToken = "hello";
broadcast(publicToken); // sink
}
```

- Line 2 reaches Line 6 along the ICFG (control-dependence holds),
- secretToken and publicToken are not aliases (data-dependence does not hold),
- secretToken@Line 2 does not flow to b@Line 6.

#### Example 3

```
char* foo(char* token){ return token: }
   int main(){
        if(condition){
3
            char* secretToken = tgetstr(...); // source
            char* b = foo(secretToken);
5
        else{
            char* publicToken = "hello";
            char* a = foo(publicToken);
            broadcast(a):
                                                 // sink
10
11
12
```

#### Example 3

```
char* foo(char* token){ return token: }
    int main(){
        if(condition){
3
            char* secretToken = tgetstr(...); // source
            char* b = foo(secretToken);
5
        elsef
            char* publicToken = "hello";
            char* a = foo(publicToken);
9
            broadcast(a):
                                                 // sink
10
11
12
```

- secretToken and a are aliases due to callee foo (data-dependence holds),
- Line 4 does not reach Line 10 on ICFG (control-dependence does not hold),
- secretToken@Line 4 does not flow to b@Line 10.

#### Example 4

```
int main(){
        char* secretToken = tgetstr(...);
                                                             // source
        while(loopCondition){
3
            if(BranchCondition){
                char* a = secretToken;
                broadcast(a);
                                                           // sink
            else{
                char* b = "hello";
10
11
12
```

How many tainted flows from source to sink?

#### Example 4

```
int main(){
        char* secretToken = tgetstr(...);
                                                             // SOUTCE
        while(loopCondition){
            if (BranchCondition) {
                 char* a = secretToken:
                 broadcast(a);
                                                           // sink
            elsef
                 char* b = "hello";
10
11
12
```

How many tainted flows from source to sink?

- Two paths from Line 2 to Line 6 along the ICFG (control-dependence holds),
- secretToken and a are aliases (data-dependence holds),
- secretToken@Line 2 have two tainted paths flowing to b@Line 6.

## **Configuring Sources and Sinks for Taint Analysis**

**Aim**: enable different taint tracking patterns by defining/configuring sources and sinks.

 Given a source v<sub>src</sub>@s<sub>src</sub> and a sink v<sub>snk</sub>@s<sub>snk</sub>, in this class, we are interested in the case that s<sub>src</sub> and s<sub>snk</sub> are both API calls, i.e., CallBlockNode in SVF.

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- v<sub>src</sub> is a return value from the call statement s<sub>src</sub>.
- $\mathbf{v}_{snk}$  is a parameter passing to a call statement  $\mathbf{s}_{snk}$ .

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- v<sub>src</sub> is a return value from the call statement s<sub>src</sub>.
- $\mathbf{v}_{snk}$  is a parameter passing to a call statement  $\mathbf{s}_{snk}$ .
- We can identify s<sub>src</sub> and s<sub>snk</sub> according to different APIs, so as to configure sources and sinks.
- In our Example 1, variable secretToken is V<sub>src</sub> and b is V<sub>snk</sub>. The call statement tgetstr(..) represents S<sub>src</sub> and broadcast(..) are used for S<sub>snk</sub>.

- Code template and specification: https://github.com/SVF-tools/SVF-Teaching/wiki/Assignment-4
- Make sure your previously implementations in Assignment-2.cpp and Assignment-3.cpp are in place.
  - Class TaintGraphTraversal in Assignment 4 is a child class of 'ICFGTraversal'. TaintGraphTraversal will use the DFS method implemented in Assignment 2 for control-flow traversal.
  - Andersen's analysis implemented in Assignment 3 will also be used for checking aliases between two pointers.

- You will need to implement tree tasks in Assignment 4 with the third one having 5 bonus points.
  - (1) Implement method readSrcSnkFormFile in Assignment-4.cpp using C++ file reading to configure sources and sinks (**Compulsory**).

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  - (2) Implement method printICFGPath to collect the tainted ICFG paths and add each path (a sequence of node IDs) as a string into std::set<std::string> paths similar to Assignment 2 (Compulsory)

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  - (1) Implement method readSrcSnkFormFile in Assignment-4.cpp using C++ file reading to configure sources and sinks (Compulsory).
  - (2) Implement method printICFGPath to collect the tainted ICFG paths and add each path (a sequence of node IDs) as a string into std::set<std::string> paths similar to Assignment 2 (Compulsory)
  - (3) Dump the taint program paths into a text file and implement a VSCode extension to annotate and visualize the tainted paths from a source to a sink (5 bonus points).

## C++ File Reading

Implement method readSrcSnkFormFile in Assignment-4.cpp to parse the two lines from SrcSnk.txt in the form of

```
line 1 contains source APIs "{ api1, api2, api3 }"
line 2 contains sink APIs "{ api1, api2, api3 }"
```

Please refer to the following links (among many others) for C++ file reading:

- https://www.tutorialspoint.com/cplusplus/cpp\_files\_streams.htm
- https://www.cplusplus.com/doc/tutorial/files/
- https://linuxhint.com/cplusplus\_read\_write/
- https://opensource.com/article/21/3/ccc-input-output

## **Visualizing Tainted Paths (5 bonus points)**

This task is optional and there is no uniform answer! Some hits as below but you are also encouraged to design and implement your own approach.

- Output taint paths into a text file in the following format for example, ' $\{$  In: number cl: number, fl: name  $\} \rightarrow \{$  In: number, cl: number, fl: name  $\} \rightarrow \{$  In: number, cl: number, fl: name  $\}$ '.
- Create a VSCode extension to read the text file
- Annotate the target source file (e.g., example.c) based on the taint paths reading from the text file.

Two VSCode extension examples (note that they are just general examples for references and are **NOT** the solution to the task):

```
https://github.com/akshatsinghkaushik/vscode-extension-example
https://github.com/spcidealacm/codepointer_js
```