



Storage &
Indexing

BITS, PILANI – K. K. BIRLA GOA CAMPUS

Database Systems

(IS F243)

by

Mrs. Shubhangi Gawali

Dept. of CS and IS





RAID Levels

(courtesy : The University of Sydney)



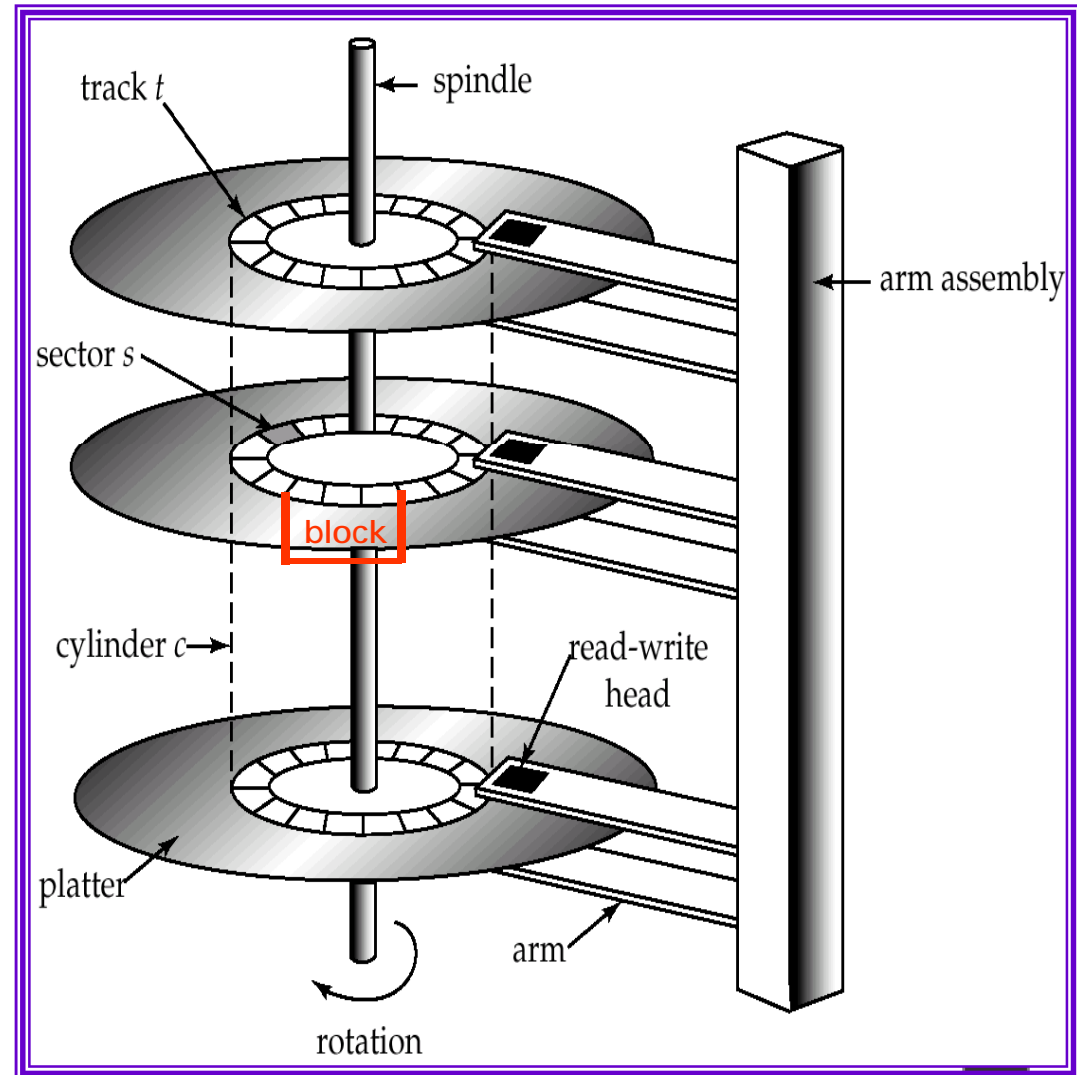
Storage & Indexing

Storage and Indexing

(courtesy : The University of Sydney)

Components of a Disk

- The platters spin (say, 120rps).
- The arm assembly is moved in or out to position a head on a desired track. Tracks under heads make a cylinder (imaginary!).
- Only one head reads/writes at any one time.
- Block size is a multiple of sector size (which is fixed).



Accessing a Disk Page

- Time to access (read/write) a disk block:
 - ▶ seek time (moving arms to position disk head on track)
 - ▶ rotational delay (waiting for block to rotate under head)
 - ▶ transfer time (actually moving data to/from disk surface)
- Seek time and rotational delay dominate.
 - ▶ Seek time varies from about 1 to 20msec
 - ▶ Rotational delay varies from 0 to 10msec
 - ▶ Transfer rate is about 1msec per 4KB page
- Key to lower I/O cost: reduce seek/rotation delays! Hardware vs. software solutions?



- Disk Array: arrangement of several disks to increase performance and improve reliability of storage system.
- RAID: Redundant Arrays of Independent Disks
 - ▶ Data striping + redundancy
- Data striping
 - ▶ distribute data over several disks
 - High capacity and high speed
 - ▶ the more disk,, the lower reliability
 - e.g., a system with 100 disks, each with MTTF of 100,000 hours (approx. 11 years), will have a system MTTF of 1000 hours (approx. 41 days)
- Redundancy
 - ▶ redundant information is maintained
 - high reliability by storing data redundantly, so that data can be recovered even if a disk fails



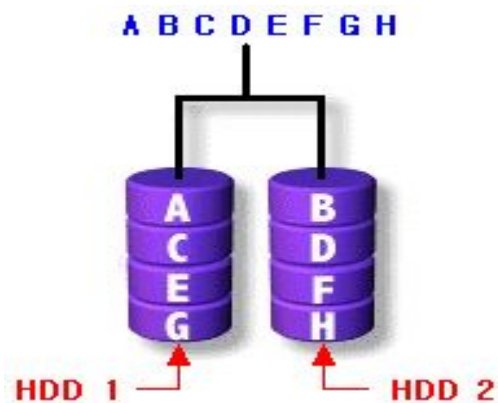
■ Storage

- ▶ Disk
- ▶ Buffer management
- ▶ File organization

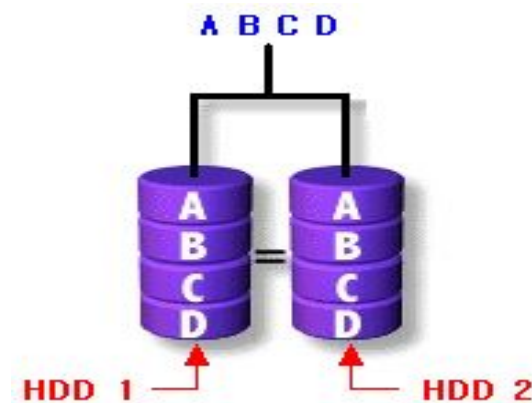
■ Indexing

- ▶ Tree-structured Indexing
- ▶ Hash-based Indexing

- Schemes to provide redundancy at lower cost by using disk striping combined with parity bits
 - ▶ Different RAID organizations, or RAID levels, have differing cost, performance and reliability characteristics
- RAID Level 0: Block striping; non-redundant.
 - ▶ Used in high-performance applications where data lost is not critical.
- RAID Level 1: Mirrored disks with block striping
 - ▶ Offers best write performance.
 - ▶ Popular for applications such as storing log files in a database system.

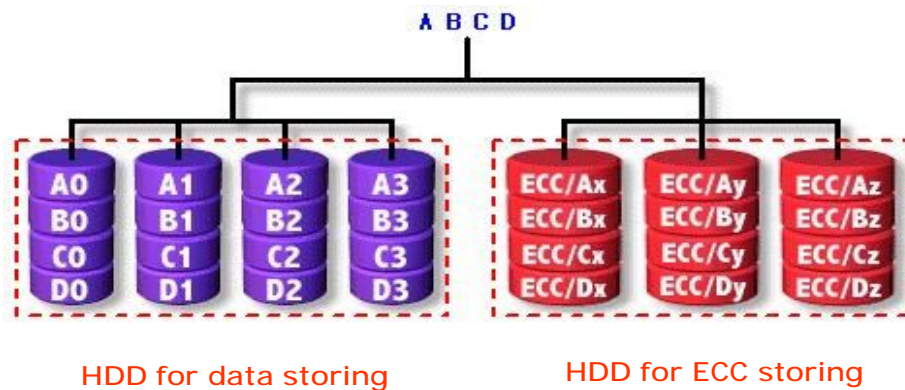
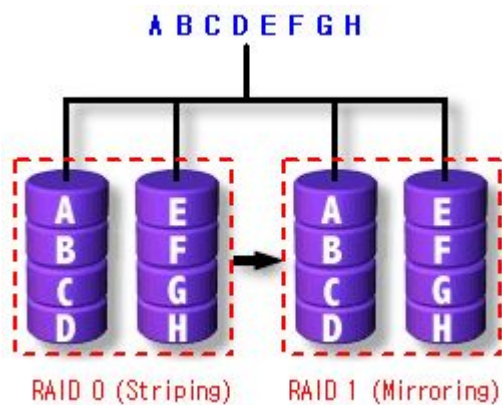


RAID 0: nonredundant striping



RAID 1: mirrored disks

- RAID Level 0+1: Striping and Mirroring
 - ▶ Parallel reads, a write involves two disks.
- RAID Level 2: Memory-Style Error-Correcting-Codes (ECC) with bit striping.
 - ▶ Striping unit is single bit
 - ▶ Store code for error correcting



RAID 0+1: striping and mirroring

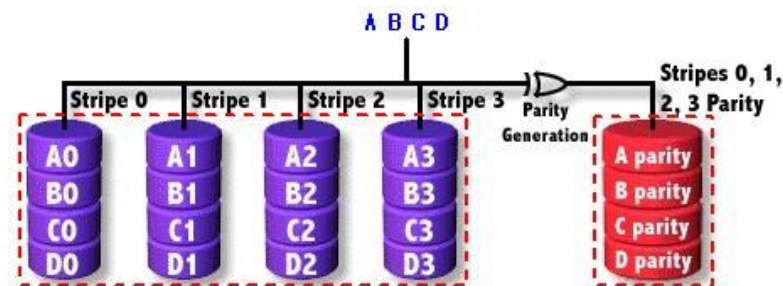
RAID 2: error correcting codes

■ RAID Level 3: Bit-Interleaved Parity

- ▶ a single parity bit is enough for error correction, since we know which disk has failed
 - When writing data, corresponding parity bits must also be computed and written to a parity bit disk

■ RAID Level 4: Block-Interleaved Parity;

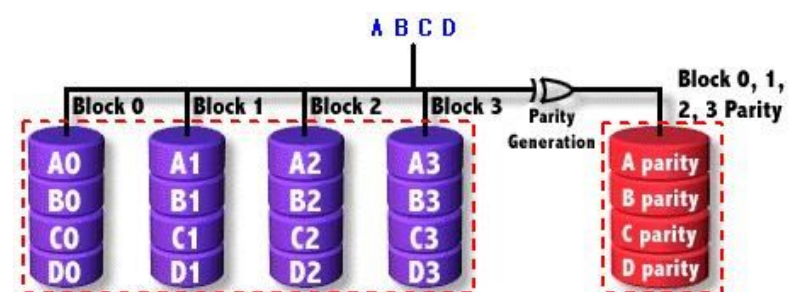
- ▶ uses block-level striping, and keeps a parity block on a separate disk for corresponding blocks from N other disks.



HDD for data storing

HDD for parity storing

RAID 3: bit-interleaved parity



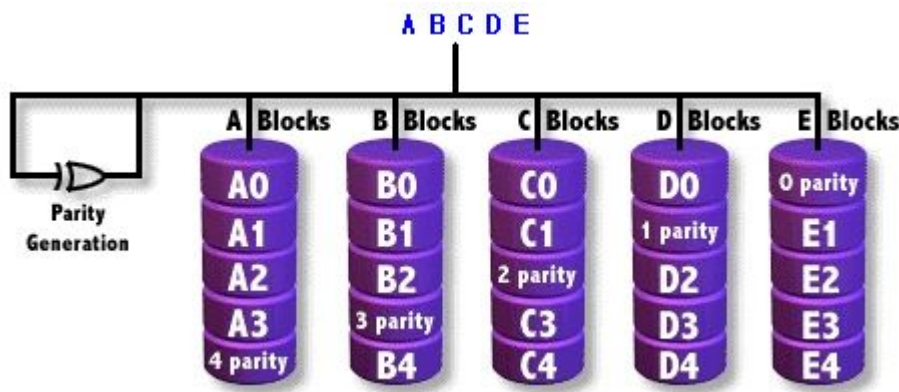
HDD for data storing

HDD for parity storing

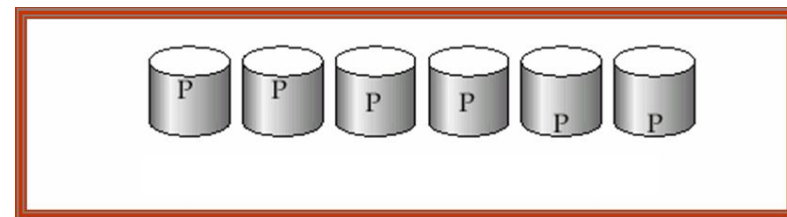
RAID 4: block-interleaved parity

RAID Levels (Cont.)

- RAID Level 5: Block-Interleaved Distributed Parity;
 - ▶ partitions data and parity among all $N + 1$ disks, rather than storing data in N disks and parity in 1 disk.
 - E.g., with 5 disks, parity block for n th set of blocks is stored on disk $(n \bmod 5) + 1$, with the data blocks stored on the other 4 disks.
- RAID Level 6: P+Q Redundancy scheme; similar to Level 5, but stores extra redundant information to guard against multiple disk failures.
 - ▶ Better reliability than Level 5 at a higher cost; not used as widely.



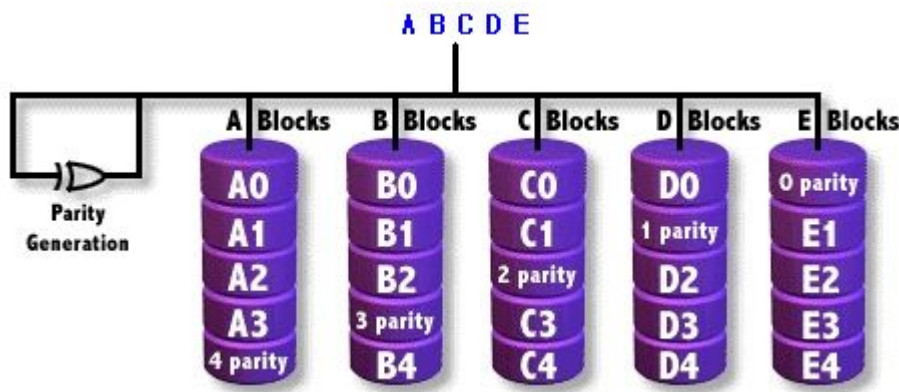
RAID 5: block-interleaved distribute parity



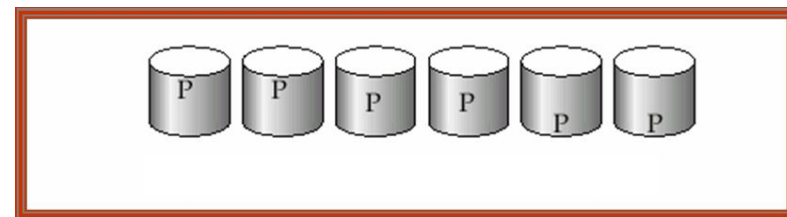
RAID 6: P+Q redundancy schem

Example of RAID Levels

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- RAID Level 6: P+Q Redundancy scheme; similar to Level 5, but stores extra redundant information to guard against multiple disk failures.
 - ▶ Better reliability than Level 5 at a higher cost; not used as widely.



RAID 5: block-interleaved distribute parity



RAID 6: P+Q redundancy schem



Choice of RAID Level

- Factors in choosing RAID level
 - ▶ Monetary cost
 - ▶ Performance: # of I/Os per second and bandwidth during normal operation
 - ▶ Performance during failure
 - ▶ Performance during rebuild of failed disk / time to rebuild failed disk
- RAID 0 is used only when data safety is not important
 - ▶ e.g. data can be recovered quickly from other sources
- Level 2 and 4 never used since they are subsumed by 3 and 5
- Level 3 is not used anymore since bit-striping forces single block reads to access all disks, wasting disk arm movement, which block striping (level 5) avoids
- Level 6 is rarely used since levels 1 and 5 offer adequate safety for almost all applications
- So competition is between 1 and 5 only
 - ▶ Level 5 is preferred for applications with low update rate, and large amounts of data
 - ▶ Level 1 is preferred for all other applications