



**BITS Pilani**  
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# Operating Systems

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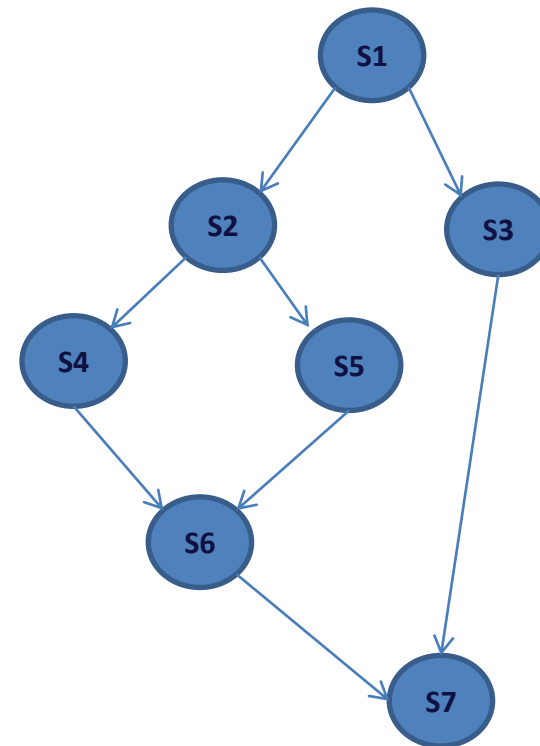
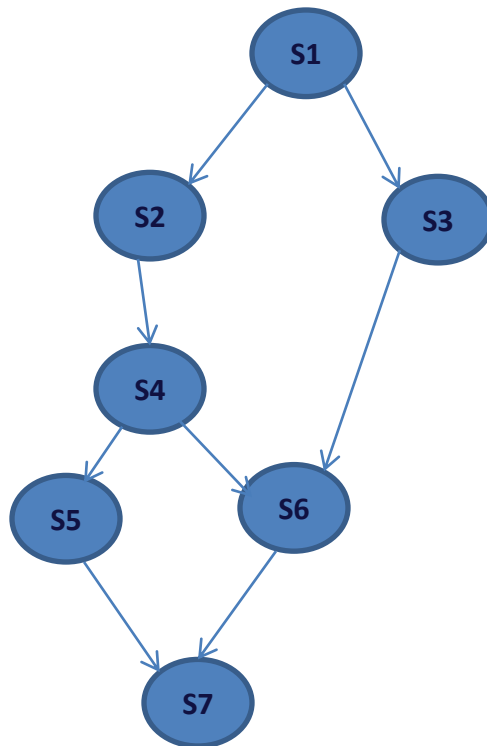


# Process Synchronization

# Last class



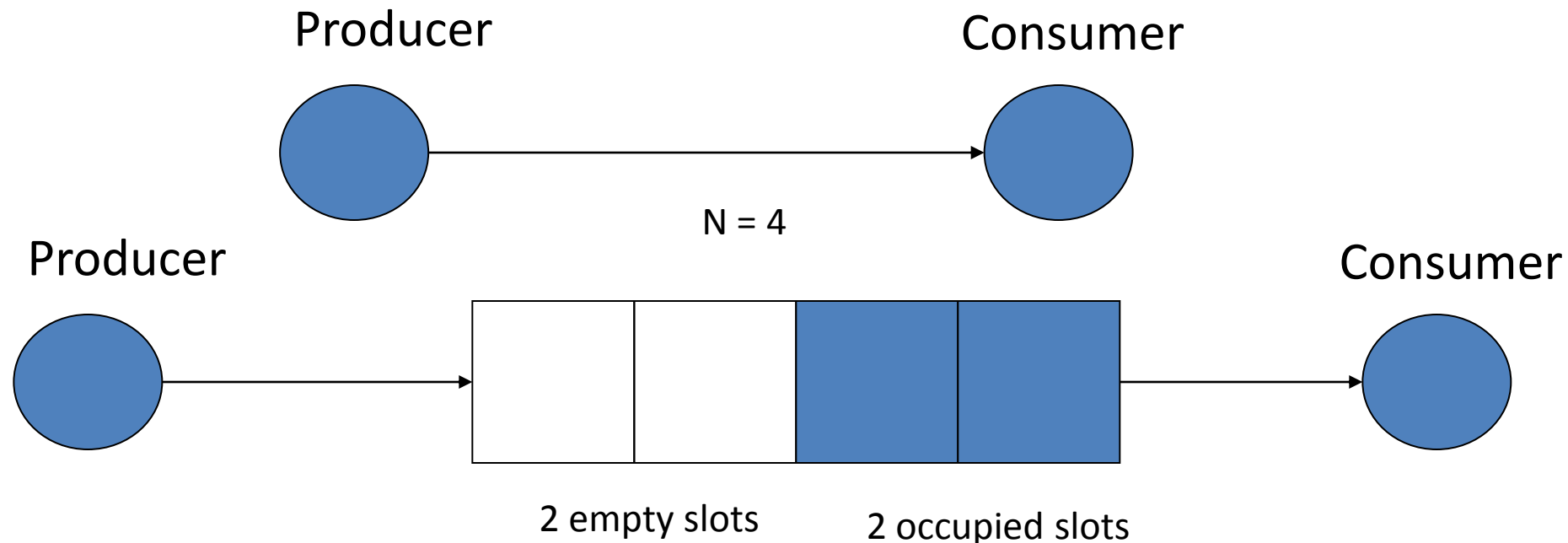
- Concurrent Processing
- fork and join construct
- cobegin and coend



# contd...



- Concurrent access to shared data may result in data inconsistency
- Maintaining data consistency requires mechanisms to ensure the orderly execution of cooperating processes
- Producer / Consumer Problem



# Producer / Consumer Problem...



- Two types of buffers:
  - unbounded buffer:
    - No limit on the size of the buffer
    - always space for producer to store data items
    - if producer is slower consumer needs to wait
  - bounded buffer: 2 cases
    - if the producer is faster than consumer
      - some point in time buffer will be full, producer has to wait
    - if the consumer is faster than producer
      - some point in time buffer will be empty, consumer has to wait
- use shared memory

- Bounded buffer can be used to enable processes to share memory.
- Shared bounded buffer is implemented using circular array.
- Assume size of buffer is 5
- Two logical pointers “in” and “out”
- variable “in” is used by producer which points to the next free position in the buffer
- variable “out” is used by consumer which points to the first full position in the buffer.

# Contd...



```
#define BUFFER_SIZE 5
int buffer [BUFFER_SIZE ];
int in = 0;
int out = 0
```

- when  $in = out$ , buffer is empty
  - when  $((in+1) \% BUFFER\_SIZE) == out$  , buffer is full
- when

# The producer process



```
while (true) {  
    int nextProduced;  
    /* produce an item and put in nextProduced */  
    while ( ( (in+1) % BUFFER_SIZE ) == out)  
        ; // do nothing and wait  
    buffer [in] = nextProduced;  
    in = (in + 1) % BUFFER_SIZE;  
}
```



# Consumer Process



```
while (true) {  
    int nextConsumed;  
    while (in == out)  
        ; // do nothing, but wait  
    nextConsumed = buffer[out];  
    out = (out + 1) % BUFFER_SIZE;  
    /* consume the item in nextConsumed  
}
```

# Producer

# Consumer



```
while (true) {  
    int nextProduced;  
    /* produce an item and put in nextProduced */  
    while ( ( (in+1) % BUFFER_SIZE ) == out )  
        ; // do nothing and wait  
    buffer[in] = nextProduced;  
    in = (in + 1) % BUFFER_SIZE;  
}
```

```
while (true) {  
    int nextConsumed;  
    while (in == out)  
        ; // do nothing, but wait  
    nextConsumed = buffer[out];  
    out = (out + 1) % BUFFER_SIZE;  
    /* consume the item in nextConsumed  
}
```

# Contd...



- Main flaw: one location has to be empty
- use counter to eliminate above problem
- Initially counter is set to zero
  - increment the counter whenever producer produces an item
  - decrement the counter whenever consumer consumes an item
  - When counter = 0 → buffer is empty
  - When counter = BUFFER\_SIZE → buffer is full

# The producer process



```
while (true) {  
    int nextProduced;  
    /* produce an item and put in nextProduced */  
    while (count == BUFFER_SIZE)  
        ; // do nothing  
    buffer [in] = nextProduced;  
    in = (in + 1) % BUFFER_SIZE;  
    count++;  
}
```

# Consumer Process



```
while (true) {  
    int nextConsumed;  
    while (count == 0)  
        ; // do nothing  
    nextConsumed = buffer[out];  
    out = (out + 1) % BUFFER_SIZE;  
    count--;  
    /* consume the item in nextConsumed  
}
```

# Producer - Consumer

## The producer process

```
while (true) {  
    int nextProduced;  
    /* produce an item and put in nextProduced */  
    while (count == BUFFER_SIZE)  
        ; // do nothing  
    buffer[in] = nextProduced;  
    in = (in + 1) % BUFFER_SIZE;  
    count++;  
}
```

## Consumer Process

```
while (true) {  
    int nextConsumed;  
    while (count == 0)  
        ; // do nothing  
    nextConsumed = buffer[out];  
    out = (out + 1) % BUFFER_SIZE;  
    count--;  
    /* consume the item in nextConsumed */  
}
```

# Contd...



## Implementation of counter

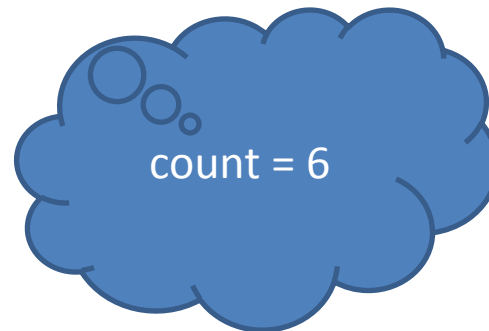
### Producer Process

A :  $R1 = \text{count}$   
B :  $R1 = R1 + 1$   
C :  $\text{count} = R1$

### Consumer Process

D :  $R2 = \text{count}$   
E :  $R2 = R2 - 1$   
F :  $\text{count} = R2$

A  
B  
D  
E  
C  
F



A :  $R1 = 6$   
B :  $R1 = 7$   
D :  $R2 = 6$   
E :  $R2 = 5$   
C :  $\text{count} = 7$   
F :  $\text{count} = 5$

# Race Condition



- Several processes access and manipulate the same data concurrently and the outcome of the execution depends on the particular order in which the access takes place, is called a **race condition**.