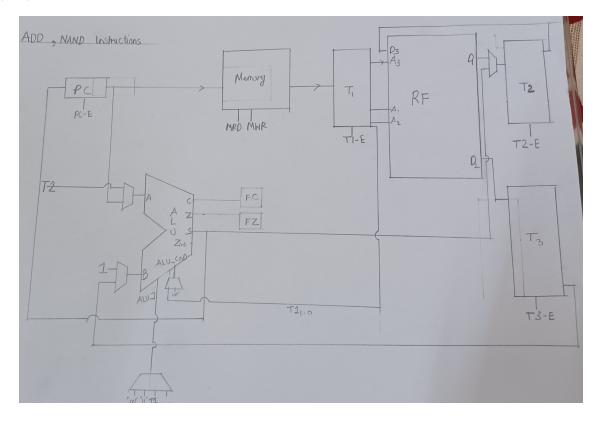
EE 224 Course Project : CPU

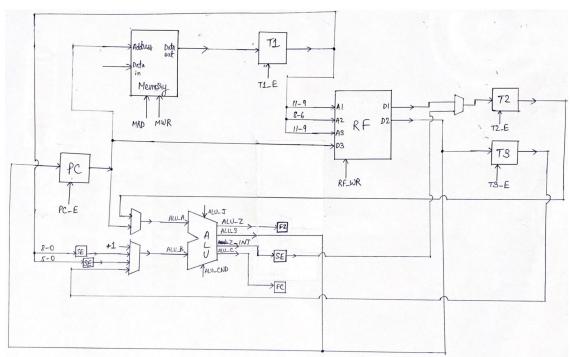
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ALU:

Data-flow





We have separated data-flow diagrams for ADD, NAND instructions from other instructions for simplicity.

Inputs:

ALU_A: Takes 3-bit input ALU_B: Takes 3-bit input

Registers storing the flags:

FC and FZ

These are connected to ALU_C and ALU_Z respectively.

Control Signals:

ALU_J: Takes 2-bit input. This specifies whether to perform addition, subtraction or NAND.

ALU_CND: Takes 2-bit input. This is used to find the new values to be updated to FZ and FC(may or may not be the same as the previous values in FZ and FC) based upon ALU_J(see table).

Outputs:

ALU_C: Outputs the carry flag to be put in the FZ register at the end of the clock cycle.

ALU_Z: Outputs the zero flag to be put in the FZ register at the end of the clock cycle. This may or may not be the same as Z_int.

ALU_S: Outputs sum, NAND, or the difference based on the bits provided in ALU_J.

Z_int: Evaluates to 1 when ALU_S is zero. Else, it evaluates to 0.

ALU_J	Function
00	Addition
01	NAND
11	Subtraction

ALU_J	ALU_CND	Output from ALU_C	Output from ALU_Z
00 (Add)	00	Modified value of carry flag.	Modified value of zero flag.
00 (Add)	10	Modified value of carry flag if input	Modified value of zero flag if input
		FC is 1. Same as the previous	FC is 1. Same as the previous
		value in FC if FC is 0.	value in FZ if FC is 0.
00 (Add)	01	Modified value of carry flag if input	Modified value of zero flag if input
		FZ is 1. Same as the previous	FZ is 1. Same as the previous
		value in FC if FZ is 0.	value in FZ if FZ is 0
00 (Add)	11 (Used for	Same as the previous value in FC.	Same as the previous value in FZ.
	updating PC)		
01 (NAND)	00	Same as the previous value in FC.	Modified value of zero flag.
01 (NAND)	10	Same as the previous value in FC.	Modified value of zero flag if input
			FC is 1. Same as the previous
			value in FZ if FC is 0.
01 (NAND)	01	Same as the previous value in FC.	Modified value of zero flag if input
			FZ is 1. Same as the previous
			value in FZ if FZ is 0.
11 (Subtract)	XX	Same as the previous value in FC.	Same as the previous value in FZ.

State Descriptions

 $(PC \equiv R7)$

 S_0 (Fetching instruction from memory)

Data Transfer	Commands
$PC \rightarrow M_{-}add$	MDR
$M_{data} \rightarrow T1$	$T1_{-}E$

 S_1 (Updating PC)

Data Transfer	Commands
$PC \rightarrow ALU_A$	PC_E
$+1 \rightarrow ALU_B$	$ALU_J \leftarrow 00$
$ALU_CND \leftarrow 11$	
$ALU_S \rightarrow PC$	

 S_2 (Reading operands)

Data Transfer	Commands
$T1_{11-9} \rightarrow RF_A1$	T2_E
$T1_{8-6} \rightarrow RF_A2$	T3_E
$RF_D1 \rightarrow T2$	
$RF_D2 \rightarrow T3$	

 S_3 (Execution)

Data Transfer	Commands
$T2 \rightarrow ALU_A$	$T2_{-}E$
$T3 \rightarrow ALU_B$	$ALU_J \leftarrow T1_{14-13}$
$T1_{1-0} \rightarrow \text{ALU_CND}$	
$ALU_S \rightarrow T2$	
$ALU_C \rightarrow FC$	
$\mathrm{ALU}_{-\!}\mathrm{Z} o \mathrm{FZ}$	

 S_4 (Storing the output)

Data Transfer	Commands
$T2 \rightarrow RF_D3$	RF_WR
$T1_{5-3} \rightarrow RF_A3$	

 S_5 (Reading operands (for ADI))

Data Transfer	Commands
$T1_{11-9} \rightarrow \text{RF-A1}$	T2_E
$RF_D1 \rightarrow T2$	T3_E
$T1_{5-0} \rightarrow \text{SE}_{-6} \rightarrow \text{T3}$	

 S_6 (Evaluating condition for BEQ)

Data Transfer	Commands
$T2 \rightarrow ALUA$	$ALU_{-}J \leftarrow 11$
$T3 \rightarrow ALU_B$	$T2_{-}E$
$Z_{int} \rightarrow SE_{1} \rightarrow T2$	
$ALU_CND \leftarrow 00$	

 S_7 (Updating PC in BEQ)

Data Transfer	Commands
$\mathrm{PC} o \mathrm{ALU}$ _A	$ALU_J \leftarrow 00$
if $(T2_0 == 0)$ then $+1 \rightarrow ALU_B$	PC_E
else $T1 \to SE_10 \to ALU_B$	
$ALU_CND \leftarrow 11$	
$ALU_S \rightarrow PC$	

 S_8 (Storing PC into REG_A)

Data Transfer	Commands
$T1_{11-9} \rightarrow RF_A3$	RF_WR
$PC \to RF_D3$	

 S_9 (Branching PC to the address PC + immediate)

Data Transfer	Commands
$\mathrm{PC} o \mathrm{ALU}$ _A	$\mathrm{ALU}_{ ext{-}\!J}$
$T1_{8-0} \rightarrow \text{SE_9} \rightarrow \text{ALU_B}$	
$ALU_CND \leftarrow 11$	
$ALU_S \rightarrow PC$	

 S_{10} (Branching PC to the address in REG_B)

Data Transfer	Commands
$T1_{8-6} \rightarrow \text{RF_A1}$	PC_E
$RF_D1 \rightarrow PC$	

 S_{11} (Executing Load Higher Immediate)

Data Transfer	Commands
$T1_{11-9} \rightarrow RF_A3$	RF_WR
$T1_{11-9} \rightarrow PZ_7 \rightarrow RF_D3$	

 S_{12} (Computing address of the memory destination)

Data Transfer	Commands
$T3 \rightarrow ALU_A$	$ALU_J \leftarrow 00$
$T1_{5-0} \rightarrow \text{SE_16} \rightarrow \text{ALU_B}$	T3_E
$ALU_S \rightarrow T3$	

 S_{13} (Writing to the memory)

Data Transfer	Commands
$T3 \rightarrow M_add$	MWR
$T2 \rightarrow M_{-}data$	

 S_{14} (Reading from memory)

Data Transfer	Commands
$T3 \rightarrow M_add$	MDR
$M_{-}data \rightarrow T2$	$T2_{-}E$

 S_{15} (Writing to the register)

Data Transfer	Commands
$T1_{11-9} \rightarrow RF_A3$	RF_WR
$T2 \rightarrow RF_D3$	

 S_{16} (Initial step of SM)

Data Transfer	Commands
$(0000000000000000) \rightarrow T2$	T2_E
$T1_{11-9} \rightarrow \text{RF_A2}$	T3_E
$RF_D2 \rightarrow T3$	

 S_{17} (Looping step 1 of SM)

Data Transfer	Commands
counter := $int(T2_{2-0})$	MWR
$T3 \rightarrow ALUA$	T3_E
$+1 \rightarrow ALU_B$	$ALU_J \leftarrow 00$
$if(T1_{counter} = = 1)$ then	
$\{T3 \rightarrow M_add$	
$T2_{2-0} \rightarrow \text{RF_A1}$	
$RF_D1 \rightarrow M_data$	
$ALU_S \rightarrow T3$ }	

 S_{18} (Looping step 2 of SM)

Data Transfer	Commands
$T2 \rightarrow ALU_A$	$ALU_J \leftarrow 00$
$1 \text{ bit} \rightarrow \text{ALU_B}$	$T2_{-}E$
$ALU_C \rightarrow T2$	

 S_{19} (Initial step of LM)

Data Transfer	Commands
$(00000000000000000) \to T2$	T2_E
$T1_{11-9} \rightarrow \text{RF_A2}$	T3_E
$RF_D3 \rightarrow T3$	

S_{20} (Looping step 1 of LM)

Data Transfer	Commands
$counter := int(T2_{2-0})$	MDR
$T1_{counter} \rightarrow RFWR$	$T3_{-}E$
$T3 \rightarrow M_add$	$ALU_J \leftarrow 00$
$M_{data} \rightarrow RF_D3$	
$T2_{2-0} \rightarrow \text{RF_A3}$	
$\mathrm{T3} ightarrow \mathrm{ALU}_{-}\mathrm{A}$	
$+1 \rightarrow ALU_B$	
$ALU_CND \leftarrow 00$	
if $(T1_{counter} = = 1)$ then ALU_S \rightarrow T3	

S_{21} (Looping step 2 of LM)

Data Transfer	Commands
$T2 \rightarrow ALU_A$	$ALU_J \leftarrow 00$
$1 \text{ bit} \rightarrow \text{ALU_B}$	$T2_E$
$ALU_C \rightarrow T2$	

Instructions with their State Diagrams and Control Signals

State flow
$S_0 \rightarrow S_1 \rightarrow S_2 \rightarrow S_3 \rightarrow S_4$
$S_0 \to S_1 \to S_2 \to S_3 \to S_4$
$S_0 \to S_1 \to S_2 \to S_3 \to S_4$
$S_0 \to S_1 \to S_5 \to S_3 \to S_4$
$S_0 \to S_1 \to S_2 \to S_3 \to S_4$
$S_0 \to S_1 \to S_2 \to S_3 \to S_4$
$S_0 o S_1 o S_2 o S_3 o S_4$
$S_0 o S_1 o S_{11}$
$S_0 \rightarrow S_1 \rightarrow S_2 \rightarrow S_{13} \rightarrow S_{15} \rightarrow S_{16}$
$S_0 \to S_1 \to S_2 \to S_{13} \to S_{14}$
$S_0 o S_1 o S_{16} o S_{17} S_{18}$ $S_0 o S_1 o S_{19} o S_{20} S_{21}$
$S_0 o S_2 o S_6 o S_7$
$S_0 ightarrow S_8 ightarrow S_9 \ S_0 ightarrow S_8 ightarrow S_{10}$

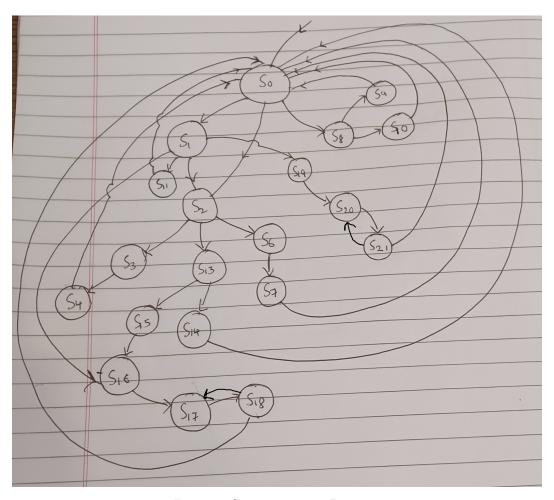


Figure 1: State Transition Diagram