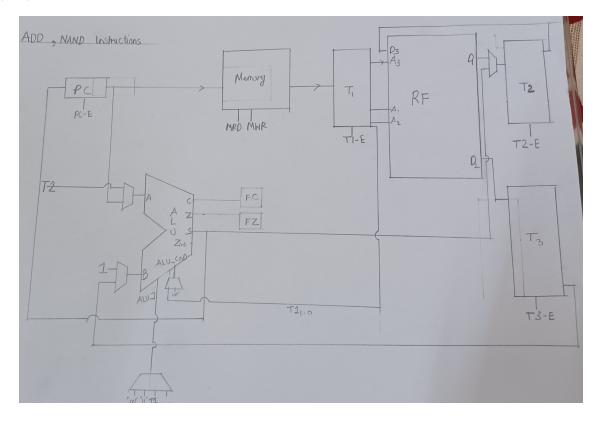
EE 224 Course Project : CPU

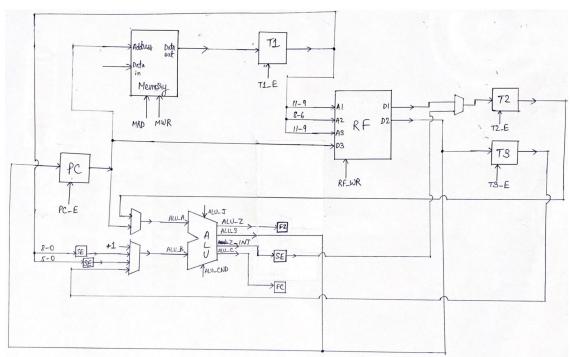
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ALU:

Data-flow





We have separated data-flow diagrams for ADD, NAND instructions from other instructions for simplicity.

Inputs:

ALU_A: Takes 16-bit input ALU_B: Takes 16-bit input

Registers storing the flags:

FC and FZ

These are connected to ALU_C and ALU_Z respectively.

Control Signals:

ALU_J: Takes 2-bit input. This specifies whether to perform addition, subtraction or NAND.

ALU_CND: Takes 2-bit input. This is used to find the new values to be updated to FZ and FC(may or may not be the same as the previous values in FZ and FC) based upon ALU_J(see table).

Outputs:

ALU_C: Outputs the carry flag to be put in the FZ register at the end of the clock cycle.

ALU_Z: Outputs the zero flag to be put in the FZ register at the end of the clock cycle. This may or may not be the same as Z_int.

ALU_S: Outputs sum, NAND, or the difference based on the bits provided in ALU_J.

Z_int: Evaluates to 1 when ALU_S is zero. Else, it evaluates to 0.

| $\mathrm{ALU}_{	extsf{-}\!J}$ | Function |
|-------------------------------|-------------|
| 00 | Addition |
| 01 | NAND |
| 11 | Subtraction |

| ALU_J | ALU_CND | Output from ALU_C | Output from ALU_Z |
|---------------|--------------|---------------------------------------|--------------------------------------|
| 00 (Add) | 00 | Modified value of carry flag. | Modified value of zero flag. |
| 00 (Add) | 10 | Modified value of carry flag if input | Modified value of zero flag if input |
| | | FC is 1. Same as the previous | FC is 1. Same as the previous |
| | | value in FC if FC is 0. | value in FZ if FC is 0. |
| 00 (Add) | 01 | Modified value of carry flag if input | Modified value of zero flag if input |
| | | FZ is 1. Same as the previous | FZ is 1. Same as the previous |
| | | value in FC if FZ is 0. | value in FZ if FZ is 0 |
| 00 (Add) | 11 (Used for | Same as the previous value in FC. | Same as the previous value in FZ. |
| | updating PC) | | |
| 01 (NAND) | 00 | Same as the previous value in FC. | Modified value of zero flag. |
| 01 (NAND) | 10 | Same as the previous value in FC. | Modified value of zero flag if input |
| | | | FC is 1. Same as the previous |
| | | | value in FZ if FC is 0. |
| 01 (NAND) | 01 | Same as the previous value in FC. | Modified value of zero flag if input |
| | | | FZ is 1. Same as the previous |
| | | | value in FZ if FZ is 0. |
| 11 (Subtract) | XX | Same as the previous value in FC. | Same as the previous value in FZ. |

State Descriptions

 $(PC \equiv R7)$

 S_0 (Fetching instruction from memory)

| Data Transfer | Commands |
|---------------------------|----------|
| $PC \rightarrow M_{-}add$ | MDR |
| $M_{data} \rightarrow T1$ | $T1_E$ |

 S_1 (Updating PC)

| Data Transfer | Commands |
|--------------------------|-----------------------|
| $PC \rightarrow ALU_A$ | PC_E |
| $+1 \rightarrow ALU_B$ | $ALU_J \leftarrow 00$ |
| $ALU_CND \leftarrow 11$ | |
| $ALU_S \rightarrow PC$ | |

 S_2 (Reading operands)

| Data Transfer | Commands |
|--------------------------------|-----------|
| $T1_{11-9} \rightarrow RF_A1$ | $T2$ _E |
| $T1_{8-6} \rightarrow RF_A2$ | $T3_{-}E$ |
| $RF_D1 \rightarrow T2$ | |
| $RF_D2 \rightarrow T3$ | |

 S_3 (Execution)

| Data Transfer | Commands |
|-----------------------------------------------|----------------------------------|
| $T2 \rightarrow ALU_A$ | $T2_{-}E$ |
| $T3 \rightarrow ALU_B$ | $ALU_{-J} \leftarrow T1_{14-13}$ |
| $T1_{1-0} \rightarrow \text{ALU_CND}$ | |
| $ALU_S \rightarrow T2$ | |
| $ALU_C \rightarrow FC$ | |
| $\mathrm{ALU}_{-\!}\mathrm{Z} 	o \mathrm{FZ}$ | |

 S_4 (Storing the output)

| Data Transfer | Commands |
|--------------------------------------|----------|
| $T2 \rightarrow RF_D3$ | RF_WR |
| $T1_{8-6} \rightarrow \text{RF_A3}$ | |

 S_5 (Reading operands (for ADI))

| Data Transfer | Commands |
|-----------------------------------------------------------|----------|
| $T1_{11-9} \rightarrow \text{RF_A1}$ | T2_E |
| $RF_D1 \rightarrow T2$ | T3_E |
| $T1_{5-0} \rightarrow \text{SE_6} \rightarrow \text{T3}$ | |

 S_6 (Evaluating condition for BEQ)

| Data Transfer | Commands |
|----------------------------------------------|--------------------------|
| $T2 \rightarrow ALUA$ | $ALU_{-}J \leftarrow 11$ |
| $T3 \rightarrow ALU_B$ | $T2_{-}E$ |
| $Z_{int} \rightarrow SE_{-1} \rightarrow T2$ | |
| $ALU_CND \leftarrow 00$ | |

 S_7 (Updating PC in BEQ)

| Data Transfer | Commands |
|--------------------------------------------------------------------|-----------------------|
| $\mathrm{PC} 	o \mathrm{ALU}$ _A | $ALU_J \leftarrow 00$ |
| $ if(T2_0 == 0) then +1 \rightarrow ALU_B$ | PC_E |
| else $T1_{5-0} \rightarrow \text{SE_6} \rightarrow \text{ALU_B}$ | |
| $ALU_CND \leftarrow 11$ | |
| $ALU_S \rightarrow PC$ | |

 S_8 (Storing PC into REG_A)

| Data Transfer | Commands |
|--------------------------------|----------|
| $T1_{11-9} \rightarrow RF_A3$ | RF_WR |
| $PC \rightarrow RF_D3$ | |

 S_9 (Branching PC to the address PC + immediate)

| Data Transfer | Commands |
|-----------------------------------------------------|--------------------------|
| $PC \to ALU_A$ | $ALU_{-}J \leftarrow 00$ |
| $T1_{8-0} \rightarrow SE_{-9} \rightarrow ALU_{-B}$ | |
| $ALU_CND \leftarrow 11$ | |
| $ALU_S \rightarrow PC$ | |

 S_{10} (Branching PC to the address in REG_B)

| Data Transfer | Commands |
|--------------------------------------|----------|
| $T1_{8-6} \rightarrow \text{RF_A1}$ | PC_E |
| $RF_D1 \rightarrow PC$ | |

 S_{11} (Executing Load Higher Immediate)

| Data Transfer | Commands |
|-----------------------------------------------|----------|
| $T1_{11-9} \rightarrow RF_A3$ | RF_WR |
| $T1_{8-0} \rightarrow PZ_7 \rightarrow RF_D3$ | |

 S_{12} (Computing address of the memory destination)

| Data Transfer | Commands |
|---------------------------------------------------------------|--------------------------|
| $T3 \rightarrow ALU_A$ | $ALU_J \leftarrow 00$ |
| $T1_{5-0} \rightarrow \text{SE_6} \rightarrow \text{ALU_B}$ | $ALU_CND \leftarrow 11$ |
| $ALU_S \rightarrow T3$ | T3_E |

 S_{13} (Writing to the memory)

| Data Transfer | Commands |
|----------------------------|----------|
| $T3 \rightarrow M_add$ | MWR |
| $T2 \rightarrow M_{-}data$ | |

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 S_{14} (Reading from memory)

| Data Transfer | Commands |
|----------------------------|----------|
| $T3 \rightarrow M_add$ | MDR |
| $M_{-}data \rightarrow T2$ | $T2$ _E |

 S_{15} (Writing to the register)

| Data Transfer | Commands |
|--------------------------------|----------|
| $T1_{11-9} \rightarrow RF_A3$ | RF_WR |
| $T2 \rightarrow RF_D3$ | |

 S_{16} (Initial step of SM)

| Data Transfer | Commands |
|---------------------------------------|----------|
| $(0000000000000000) \rightarrow T2$ | T2_E |
| $T1_{11-9} \rightarrow \text{RF_A2}$ | T3_E |
| $RF_D2 \rightarrow T3$ | |

 S_{17} (Looping step 1 of SM)

| Data Transfer | Commands |
|-----------------------------------------------------|--------------------------|
| counter := $int(T2_{2-0})$ | MWR |
| $\mathrm{T3} ightarrow \mathrm{ALU}_{-}\mathrm{A}$ | $T3_{-}E$ |
| $+1 \rightarrow ALU_B$ | $ALU_J \leftarrow 00$ |
| if $(T1_{counter} = = 1)$ then | $ALU_CND \leftarrow 11$ |
| $\{\mathrm{T3} \to \mathrm{M_add}$ | |
| $T2_{2-0} \rightarrow \text{RF_A1}$ | |
| $RF_D1 \rightarrow M_data$ | |
| $ALU_S \rightarrow T3$ } | |

 S_{18} (Updating counter variable (Looping step 2 of SM and LM))

| Data Transfer | Commands |
|-------------------------------------------|-------------------------|
| $T2 \rightarrow ALU_A$ | $ALU_J \leftarrow 00$ |
| $1 \text{ bit} \rightarrow \text{ALU_B}$ | $ALU_CND \leftarrow 11$ |
| $ALU_S \rightarrow T2$ | $T2_{-}E$ |

 S_{19} (Initial step of LM)

| Data Transfer | Commands |
|---------------------------------------|----------|
| $(0000000000000000) \rightarrow T2$ | T2_E |
| $T1_{11-9} \rightarrow \text{RF_A1}$ | T3_E |
| $RF_D1 \rightarrow T3$ | |

S_{20} (Looping step 1 of LM)

| Data Transfer | Commands |
|-------------------------------------------------------|--------------------------|
| $counter := int(T2_{2-0})$ | MDR |
| $T1_{counter} \rightarrow RF_WR$ | $T3_{-}E$ |
| $T3 \rightarrow M_add$ | $ALU_J \leftarrow 00$ |
| $M_{data} \rightarrow RF_D3$ | $ALU_CND \leftarrow 11$ |
| $T2_{2-0} \rightarrow RF_A3$ | |
| $T3 \rightarrow ALU_A$ | |
| $+1 \rightarrow ALU_B$ | |
| if $(T1_{counter} = = 1)$ then ALU_S \rightarrow T3 | |

Instructions with their State Diagrams and Control Signals

| Instruction | State flow |
|-------------|------------------------------------------------------------------------------------------------|
| ADD | $S_0 \rightarrow S_1 \rightarrow S_2 \rightarrow S_3 \rightarrow S_4$ |
| ADC | $S_0 \rightarrow S_1 \rightarrow S_2 \rightarrow S_3 \rightarrow S_4$ |
| ADZ | $S_0 \rightarrow S_1 \rightarrow S_2 \rightarrow S_3 \rightarrow S_4$ |
| ADI | $S_0 \rightarrow S_1 \rightarrow S_5 \rightarrow S_3 \rightarrow S_4$ |
| NDU | $S_0 \rightarrow S_1 \rightarrow S_2 \rightarrow S_3 \rightarrow S_4$ |
| NDC | $S_0 ightarrow S_1 ightarrow S_2 ightarrow S_3 ightarrow S_4$ |
| NDZ | $S_0 \rightarrow S_1 \rightarrow S_2 \rightarrow S_3 \rightarrow S_4$ |
| LHI | $S_0 	o S_1 	o S_{11}$ |
| LW | $S_0 \rightarrow S_1 \rightarrow S_2 \rightarrow S_{12} \rightarrow S_{14} \rightarrow S_{15}$ |
| SW | $S_0 \to S_1 \to S_2 \to S_{12} \to S_{13}$ |
| SM | $S_0 \rightarrow S_1 \rightarrow S_{16} \rightarrow S_{17} \ S_{18}$ |
| LM | $S_0 \to S_1 \to S_{19} \to S_{20} \ S_{18}$ |
| BEQ | $S_0 	o S_2 	o S_6 	o S_7$ |
| JAL | $S_0	o S_8	o S_9$ |
| JLR | $S_0 \to S_8 \to S_{10}$ |

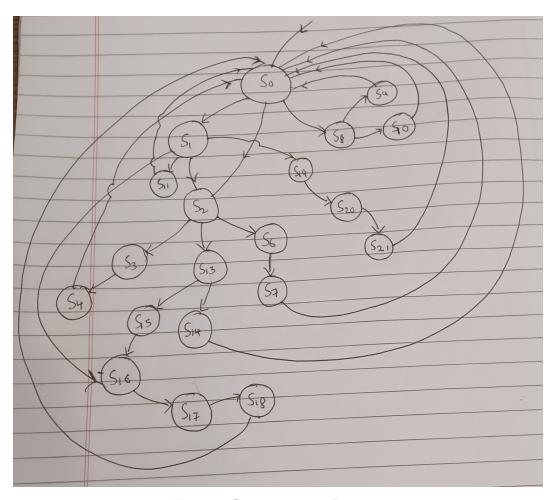


Figure 1: State Transition Diagram