

## COMP-170: Homework #3

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### Problem 4

Prove that  $L = \{\langle M \rangle \mid M \text{ accepts input } 1011\}$  is recognizable.

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We would like to show that the language  $L = \{\langle M \rangle \mid M \text{ accepts input } 1011\}$  is recognizable. In order to do this, we will use proof by construction. More specifically, we will construct a machine  $R_L$  that will accept any  $x \in L$ .

Let's define  $R_L$  as the following:

```
 $R_L$  on input  $\langle M \rangle$ :
  Run  $M$  on 1011
    If  $M$  accepts 1011, ACCEPT
    If  $M$  reject 1011, REJECT
END
```

We will now claim that  $R_L$  recognizes  $L$ . To show this, consider the following cases:

1. Suppose  $x \in L$  such that  $x = \langle M \rangle$  for some  $M$  and  $M$  accepts the input 1011. In this case, when we run  $M$  on input 1011, we know  $M$  will accept the input, which will cause  $R_L$  to accept  $x$  correctly.
2. Suppose  $y \notin L$  such that  $y = \langle M \rangle$  for some  $M$  and  $M$  rejects the input 1011. In this case, when we run  $M$  on input 1011, we know  $M$  will reject the input, which will cause  $R_L$  to reject  $y$  correctly.
3. Suppose  $z \notin L$  such that  $z = \langle M \rangle$  for some  $M$  and  $M$  loops on the input 1011. In this case, when we run  $M$  on input 1011,  $M$  will loop, which means  $R_L$  will loop. However, because  $z \notin L$ ,  $R_L$  can do anything other than accept  $z$ , which it does. Thus, we can see  $R_L$  functions correctly on  $z$ .

Thus, since  $R_L$  accepts all  $x \in L$  and doesn't accept all  $x \notin L$ , we can see that  $R_L$  recognizes  $L$ . This shows that  $L$  is indeed recognizable.  $\square$