

Dynamic Programming

- 1) Greedy Method → decision is taken one time
- 2) Dynamic Programming → decision taken at each step
↑
principle of optimality

ex:

$$\text{fib}(n) = \begin{cases} 0 & \text{if } n = 0 \\ 1 & \text{if } n = 1 \\ \text{fib}(n-2) + \text{fib}(n-1) & \text{if } n > 1 \end{cases}$$

int fib(int n)

{

if (n <= 1)

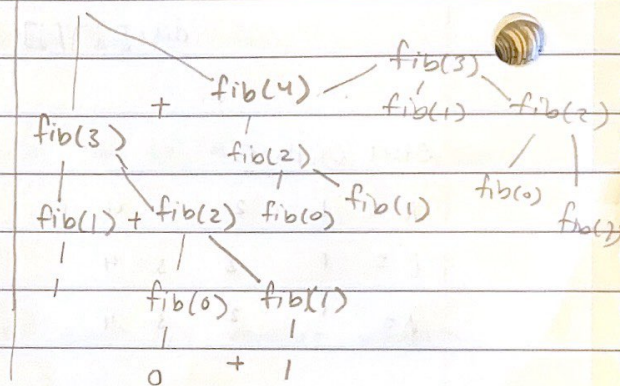
return n;

return fib(n-2) + fib(n-1);

}

0, 1, 1, 2, 3, 5, 8, 13, 21

fib(5)



$$T(N) = 2T(N-1) + 1$$

$$O(2^N)$$

Reduce function calls of fib(1) & fib(2). Since they keep getting called increasing our big O.

F	0	1	-1	-1	-1	-1	
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0 1 2 3 4 5

Start out at (-1) for all keep going through the $\text{fib}(N)$ graph solving all of these to save time complexity when solving.

F	0	1	1	2	3	5	
	0	1	2	3	4	5	

memoization has made it from $O(2^n)$ to $O(N)$, 15 calls to just 6.

memoization follows top down approach.

S 5

A 3

B 4

C 2

D 6

G 0

$$(1) 8 \rightarrow A \rightarrow B = 2 + 4 = 6$$

$$8 \rightarrow A \rightarrow C = 2 + 2 = 4$$

$$(2) 5 \rightarrow A \rightarrow B \rightarrow C = 2 + 4 + 2 = 8$$

Tabulation

F	0	1	1	2	3	5
	0	1	2	3	4	5

$$\text{fib}(n) = \begin{cases} 0 & \text{if } n=0 \\ 1 & \text{if } n=1 \\ \text{fib}(n-2) + \text{fib}(n-1) & \text{if } n>1 \end{cases}$$

```
int fib (int n)
```

```
{
```

```
    if (n <= 1)
```

```
        return n;
```

```
    F[0] = 0; F[1] = 1;
```

```
    for (int i = 2; i <= n; i++)
```

```
    {
```

```
        F[i] = F[i-2] + F[i-1];
```

```
    }
```

```
    return F[N];
```

```
}
```

fib(5)? = 5 ✓

Table is generated. Using
bottom up approach.

$T(N) = 2T(N-1) + 1$

$O(2^N)$

[Review Slide Notes for Dynamic Programming!]

BOTH Tabulation and Memoization can be used, but most
will be solved using Tabulation!