Ограничения MapReduce Spark RDI Spark Program Implementation

Apache Spark

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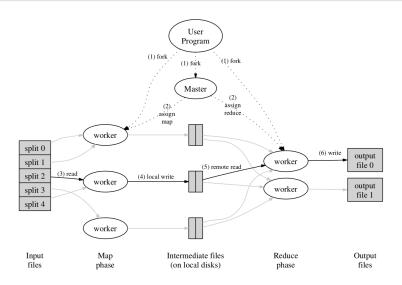
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Ограничения MapReduce Spark RDD

Spark RDD
Spark Programs
Implementation

Reminder

роблемы



Reminder Проблемы Частичные решени

- Iterative ML algorithms
- Ad-hoc analytics
- Interactive data-mining

Map Reduce is Good Enough?¹

¹Jimmy Lin. "Mapreduce is good enough? if all you have is a hammer, throw away everything that's not a nail!" B: *Big Data* 1.1 (2013), c. 28—37.

леппиет Проблемы Пастичные решения

- Hive² data wearehousing solution
- Pig³ dataflow system

²Ashish Thusoo и др. "Hive: a warehousing solution over a map-reduce framework".

B: Proceedings of the VLDB Endowment 2.2 (2009), c. 1626—1629.

³Alan F Gates и др. "Building a high-level dataflow system on top of Map-Reduce: the Pig experience". B: *Proceedings of the VLDB Endowment* 2.2 (2009), c. 1414—1425.

Hive Main components

- HiveQL SQL like language
- Metastore catalog with metadata about tables
- Compiler converts query to a execution plan (MapReduce jobs)

 Like Hive but with different query language and without Metastore

D abstraction

Пример

Linage graph / Lazy computation

Определение

RDD:

- Resilient отказоустойчивый
- Distributed разбитый на партиции
- Dataset

read-only, partitioned collection of records

Efficient Fault-tolerance

- Запомним граф вычислений (linage)
- Тогда если часть данных будет потеряна, то их легко можно восстановить

RDD creation

RDD можно построить одним из следующих способов:

- Из данных находящихся на HDFS или в RAM
- Выполнив операцию над существующим RDD:
 - Transformations
 - Actions

Implementat

DD abstraction

Пример

Linage graph / Lazy computation

Transformations

- map
- filter
- join
- reduceByKey
- ..

Actions

- count количество элементов в RDD
- save сохранение RDD, например, на HDFS
- collect отправить RDD на driver

Persistance and Partitioning

- Users can indicate which RDDs they will reuse and choose a storage strategy for them (e.g., in-memory storage).
- They can also ask that an RDD's elements be partitioned across machines based on a key in each record.

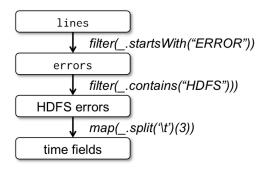
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DD abstraction ример nage graph / Lazy computatio

Операции задают декларативное описание того, что мы хотим сделать

```
1 lines = spark.textFile("hdfs://...")
2 errors = lines.filter( .startsWith("ERROR"))
  errors.persist()
  errors.count()
6
  // Count errors mentioning MvSOL:
  errors.filter( .contains("MySQL")).count()
9
  // Return the time fields of errors mentioning
11 // HDFS as an array (assuming time is field
  // number 3 in a tab—separated format):
13 errors filter( contains("HDFS"))
         .map(\_.split('\t')(3))
14
         .collect()
15
```

Lineage graph



Puc.: Boxes represent RDDs and arrows represent transformations

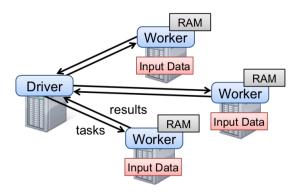
RDD abstraction Пример Linage graph / Lazy computatio

Linage graph

RDD has enough information about how it was derived from other datasets (its lineage) to compute its partitions from data in stable storage

Lazy computation

- Spark computes RDDs lazily the first time they are used in an action, so that it can pipeline transformations.
- Spark keeps persistent RDDs in memory by default, but it can spill them to disk if there is not enough RAM.



Puc.: The user's *driver* program launches multiple *workers*, which read data blocks from a distributed file system and can persist computed RDD partitions in memory.

Transformations Types

```
map(f: T \Rightarrow U):
                                          RDD[T] \Rightarrow RDD[U]
             filter(f: T \Rightarrow Bool):
                                         RDD[T] \Rightarrow RDD[T]
       flatMap(f: T \Rightarrow Sea[U]):
                                         RDD[T] \Rightarrow RDD[U]
      sample(fraction: Float) :
                                         RDD[T] \Rightarrow RDD[T]
                  groupBvKev():
                                          RDD[(K, V)] \Rightarrow RDD[(K, Seq[V])]
  reduceBvKev(f: (V, V) \Rightarrow V)
                                         RDD[(K, V)] \Rightarrow RDD[(K, V)]
                          union() : (RDD[T], RDD[T]) \Rightarrow RDD[T]
                            join()
                                          (RDD[(K, V)], RDD[(K, W)])
                                            \Rightarrow RDD[(K,(V,W))]
                                          (RDD[(K, V)], RDD[(K, W)])
                       cogroup()
                                            \Rightarrow RDD[(K, (Seq[V], Seq[W]))]
                 crossProduct() : (RDD[T], RDD[U]) \Rightarrow RDD[(T, U)]
         mapValues(f: V \Rightarrow W):
                                         RDD[(K, V)] \Rightarrow RDD[(K, W)]
       sort(c: Comparator[K]):
                                          RDD[(K, V)] \Rightarrow RDD[(K, V)]
partitionBv(p: Partitioner[K])
                                          RDD[(K, V)] \Rightarrow RDD[(K, V)]
```

Actions Types

count() : $RDD[T] \Rightarrow Long$

 $\begin{array}{ccc} collect() & : & RDD[T] \Rightarrow Seq[T] \\ reduce(f:(T,T)\Rightarrow T) & : & RDD[T]\Rightarrow T \\ lookup(k:K) & : & RDD[(K,V)]\Rightarrow Seq[V] \end{array}$

save(path: String) : Outputs RDD to a storage system

Logistic Regression

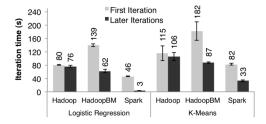
LogReg⁴

⁴Trevor Hastie, Robert Tibshirani и Jerome Friedman. *The elements of statistical learning: data mining, inference, and prediction*. Springer Science & Business Media, 2009.

Список литературы

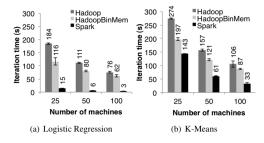
Logistic Regression

Logistic Regression Performance



Puc.: Duration of the first and later iterations in Hadoop, HadoopBinMem and Spark for logistic regression and k-means using 100 GB of data on a 100-node cluster.

Logistic Regression Performance



 ${\sf Puc.:}$ Running times for iterations after the first in Hadoop, HadoopBinMem, and ${\sf Spark}$

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Operations Примеры Representing RDD

Logistic Regression Performance

Keeping points in memory across iterations can yield a $20\times$ speedup

Operations Примеры Representing RDDs

PageRank

PageRank⁵⁶

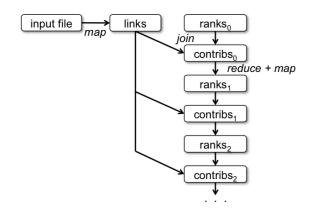
⁵Lawrence Page и др. *The pagerank citation ranking: Bringing order to the web.* Tex. отч. Stanford InfoLab, 1999.

 $^{^6}$ Jure Leskovec, Anand Rajaraman μ Jeffrey David Ullman. *Mining of massive data sets.* Cambridge university press, 2019.

PageRank Code

```
val links = spark.textFile(...).map(...).persist()
var ranks = // RDD of (URL, rank) pairs
  for (i <- 1 to ITERATIONS) {
    // Build an RDD of (targetURL, float) pairs
    // with the contributions sent by each page
    val contribs = links.join(ranks).flatMap {
       (url, (links, rank)) =>
        links.map(dest => (dest, rank/links.size))
8
9
    // Sum contributions by URL and get new ranks
10
    ranks = contribs.reduceByKey((x,y) => x+y)
11
                .mapValues(sum => a/N + (1-a)*sum)
13 }
```

PageRank Linage graph



PageRank Performance

Preserving partitioning might help

If ranks and links are co-partitioned then join requires no communication

PageRank Performance

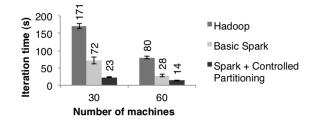
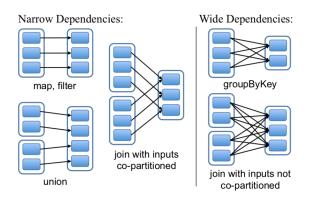


Рис.: Performance of PageRank on Hadoop and Spark.

- Partititions atomic pices of the dataset
- Dependencies dependencies on parent RDDs

Dependencies



Puc.: Examples of narrow and wide dependencies. Each box is an RDD, with partitions shown as shaded rectangles.

Dependencies Narrow

- Narrow dependencies allow for pipelined execution on one cluster node, which can compute all the parent partitions.
- Recovery after a node failure is more efficient with a narrow dependency, as only the lost parent partitions need to be recomputed, and they can be recomputed in parallel on different nodes.

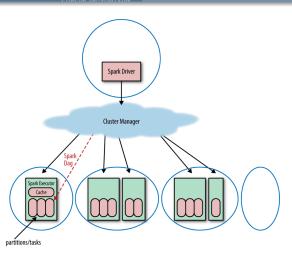


Рис.: Starting a Spark application on a distributed system

Spark Application

Замечание

- One node can have multiple Spark executors, but an executor cannot span multiple nodes.
- An RDD will be evaluated across the executors in partitions (shown as red rectangles).
- Each executor can have multiple partitions, but a partition cannot be spread across multiple executors.

SparkContext

SparkContext - connection between user's program and cluster. Containes information about requested resources, type of resources allocation (dynamic/static), etc

Pipeline

- User runs an action on RDD
- Scheduler builds a DAG of stages (each stage containes piplened transformations with narrow dependencies. Boundaries - wide deps.)
- Scheduler launches tasks to compute missing partitions

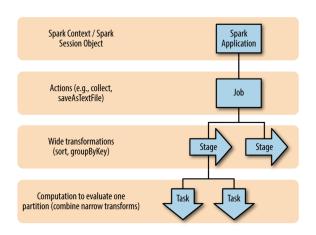
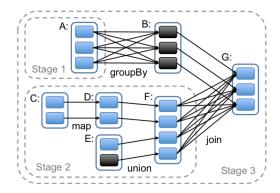


Рис.: The Spark application tree



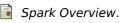
Puc.: Boxes with solid outlines are RDDs. Partitions are shaded rectangles, in black if they are already in memory.

- User can choose how to persist data:
 - in-memory storage as deserialized Java objects
 - in-memory storage as serialized data
 - on-disk storage.
- LRU eviction policy at the level of RDDs

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