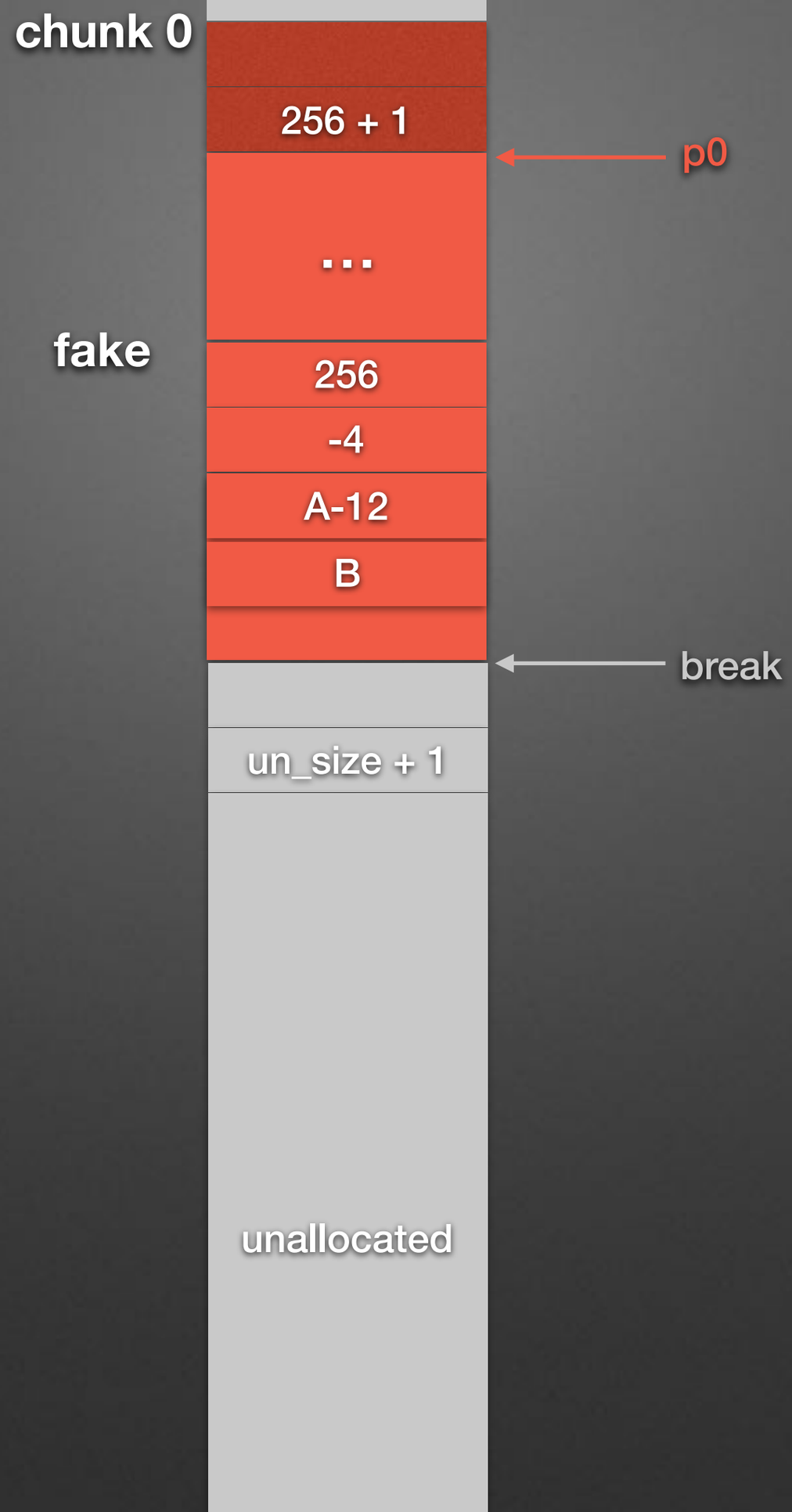
任意地址任意写

首先申请1个堆块p0，并伪造一个与p0相邻的空闲堆块fake，
在fake->presize填入偶数，fake->size填入-4，
fake->fd填入A-12，fake->bk填入B，
当释放p0时触发unlink(fake)过程，就能实现在任意地址A写入任意值B

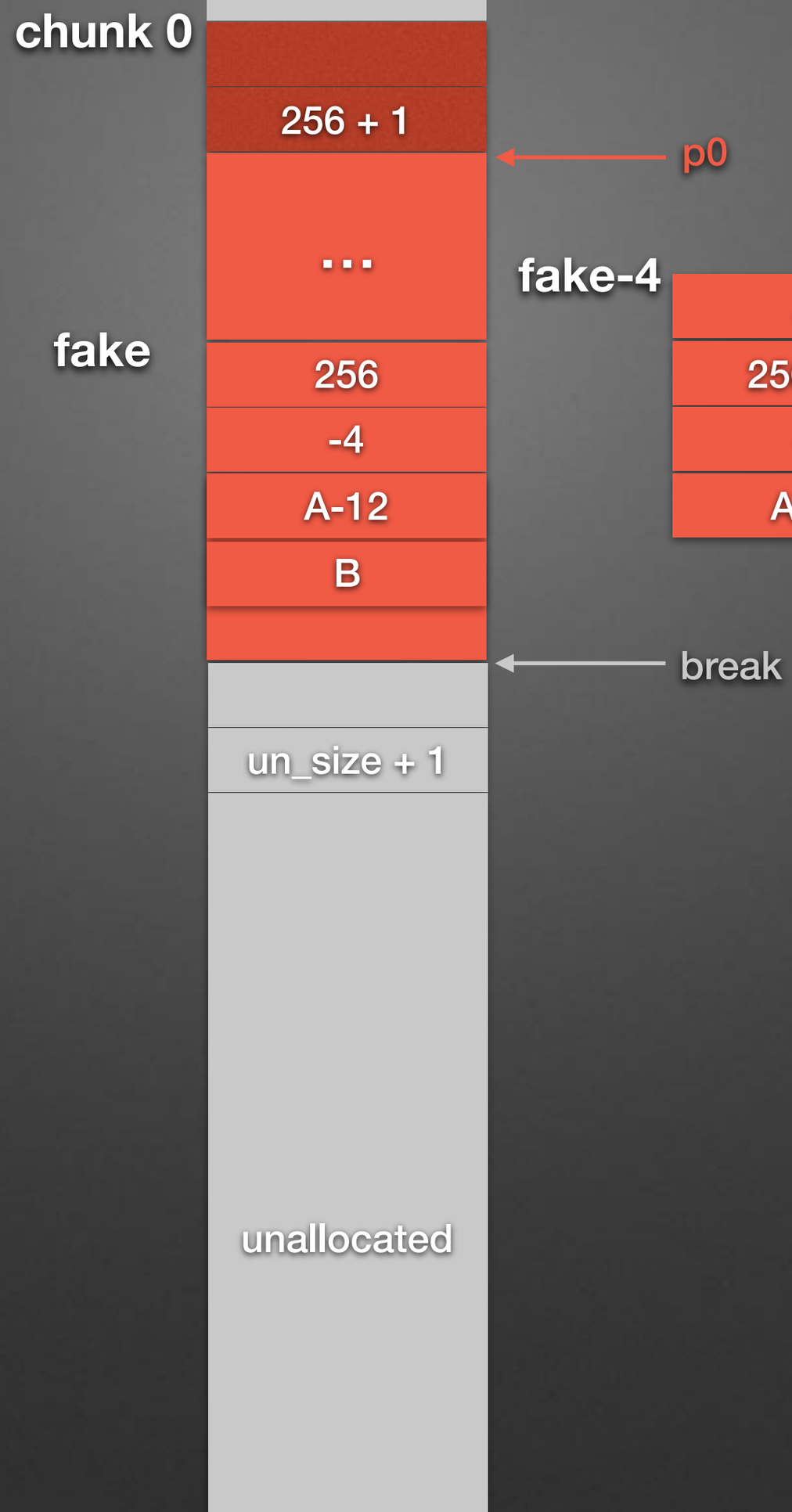
free(p0)

由chunk0 + 256找到chunk0下一个堆块fake，通过查看fake状态判断是否进行合并

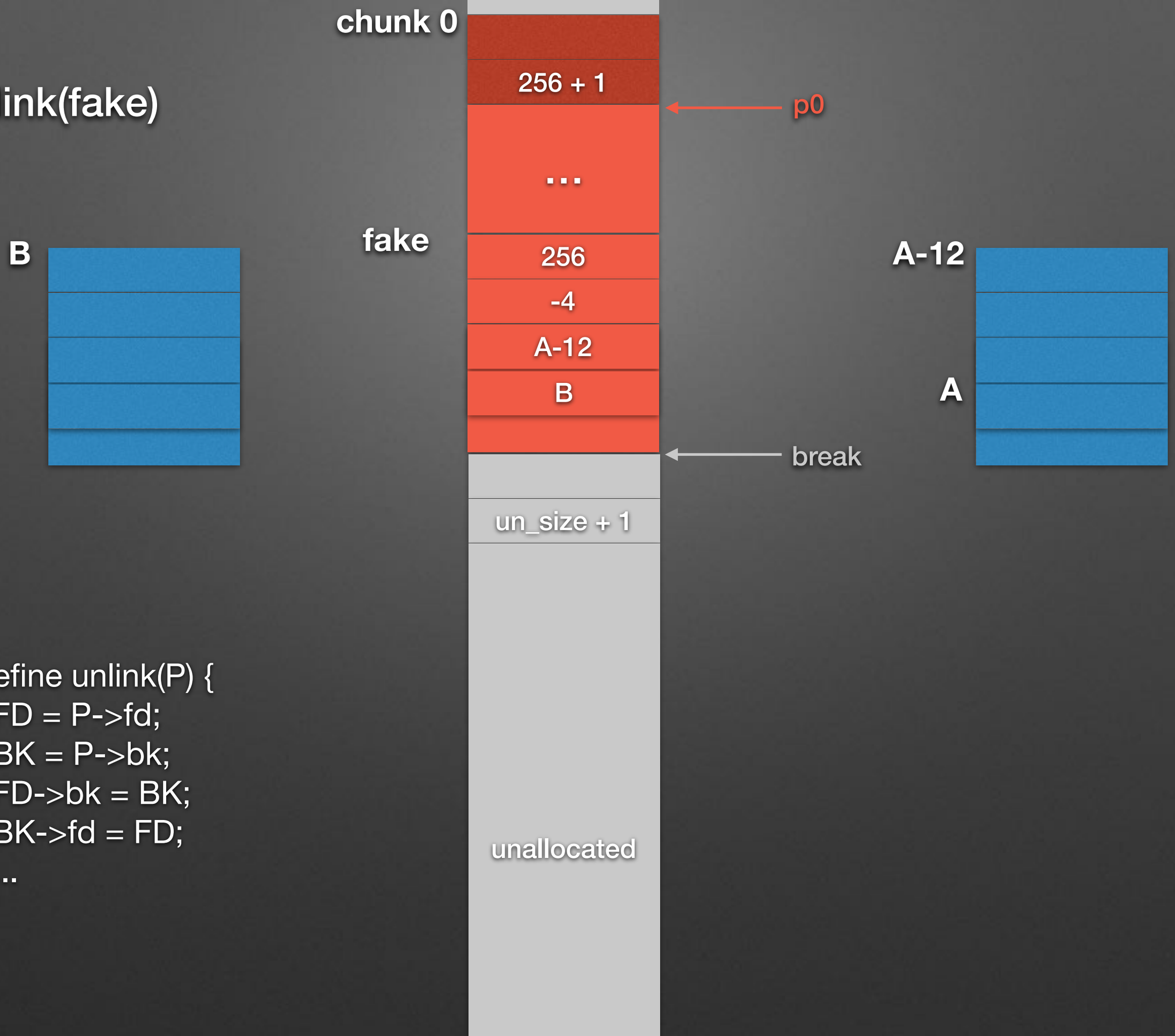


free(p0)

由于fake的size为-4,堆管理器将fake-4视为fake的下一个相邻堆块，检查发现其size的标记位为0，于是认为fake为空闲堆块，触发空闲堆块合并

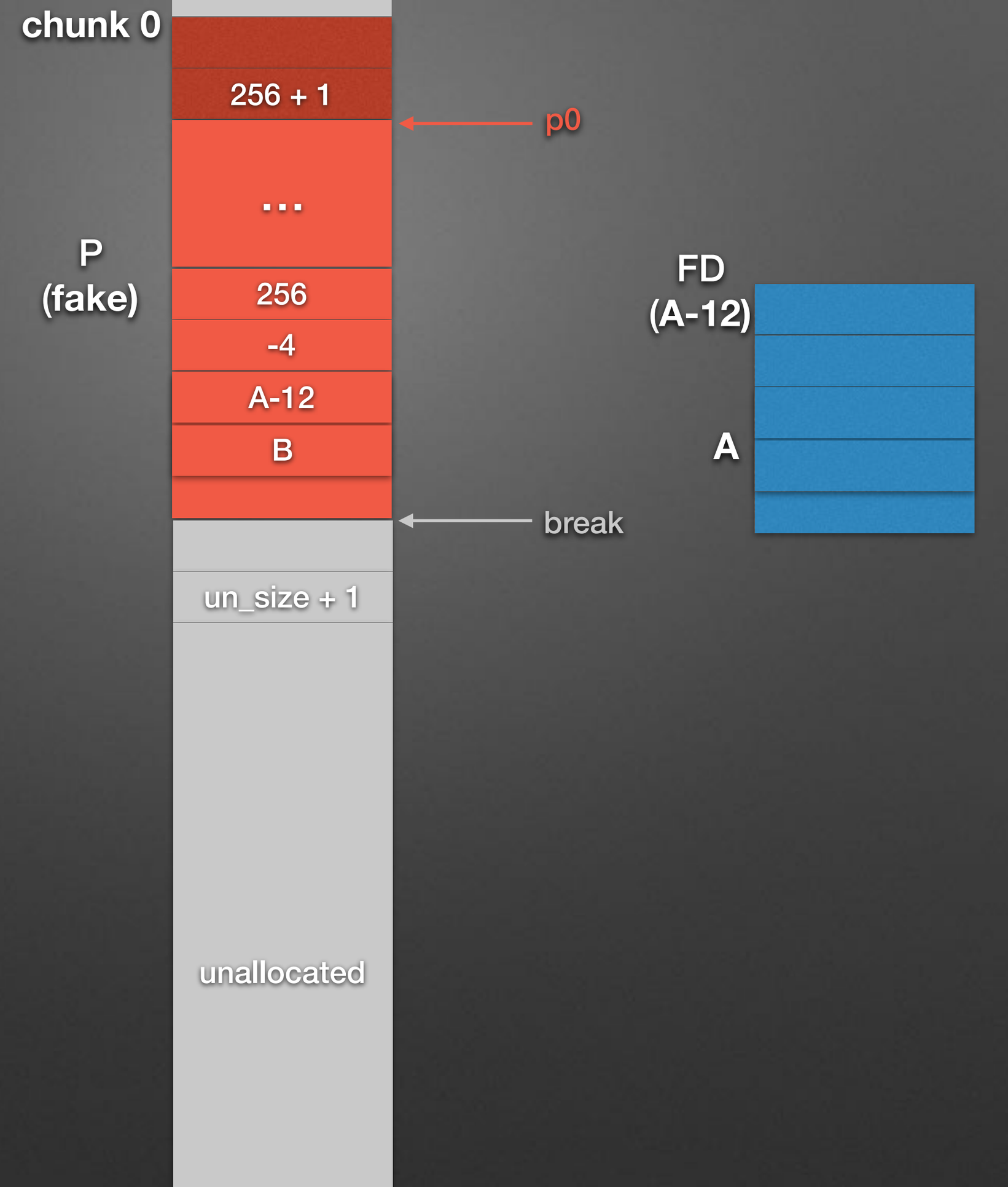


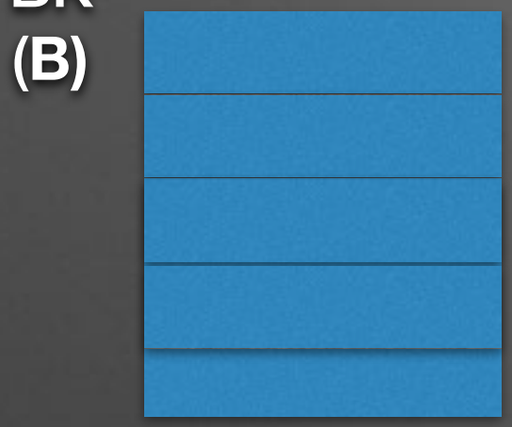
unlink(fake)




```
#define unlink(P) {
    FD = P->fd;
    BK = P->bk;
    FD->bk = BK;
    BK->fd = FD;
    ...
}
```

unlink(fake)



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#define unlink(P) {
    FD = P->fd;
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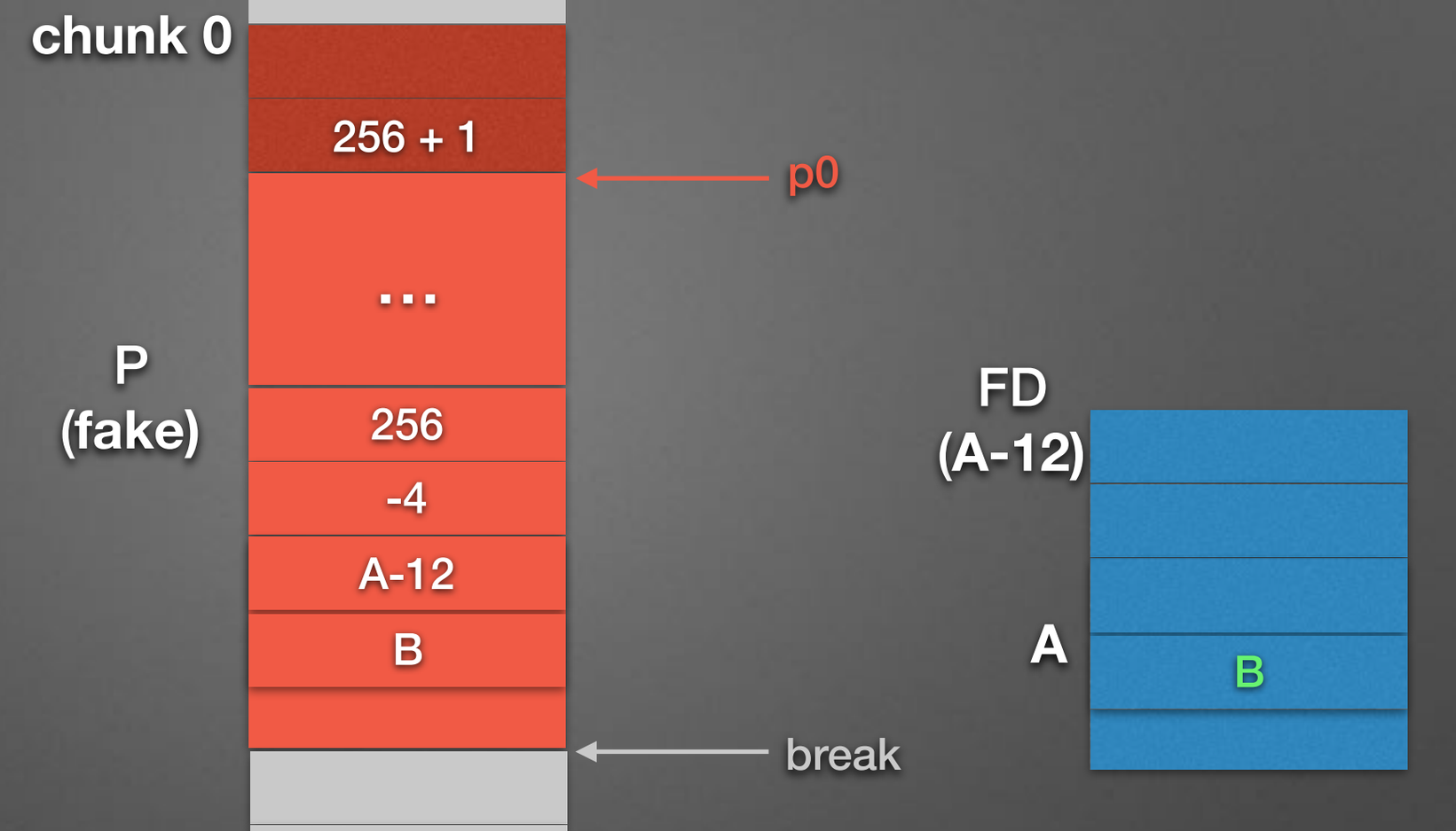
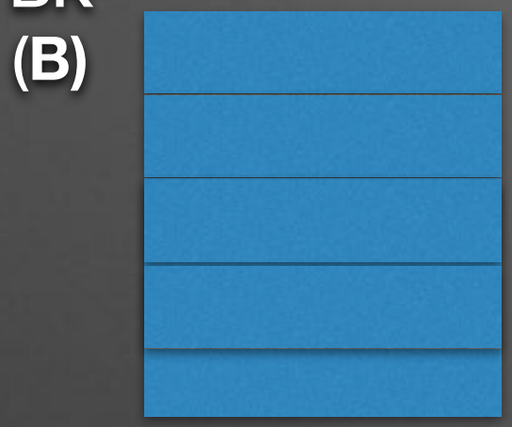
    FD->bk = BK;
    BK->fd = FD;
    ...
}
```

unlink(fake)



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#define unlink(P) {
    FD = P->fd;
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    FD->bk = BK;
    BK->fd = FD;
    ...
}
```
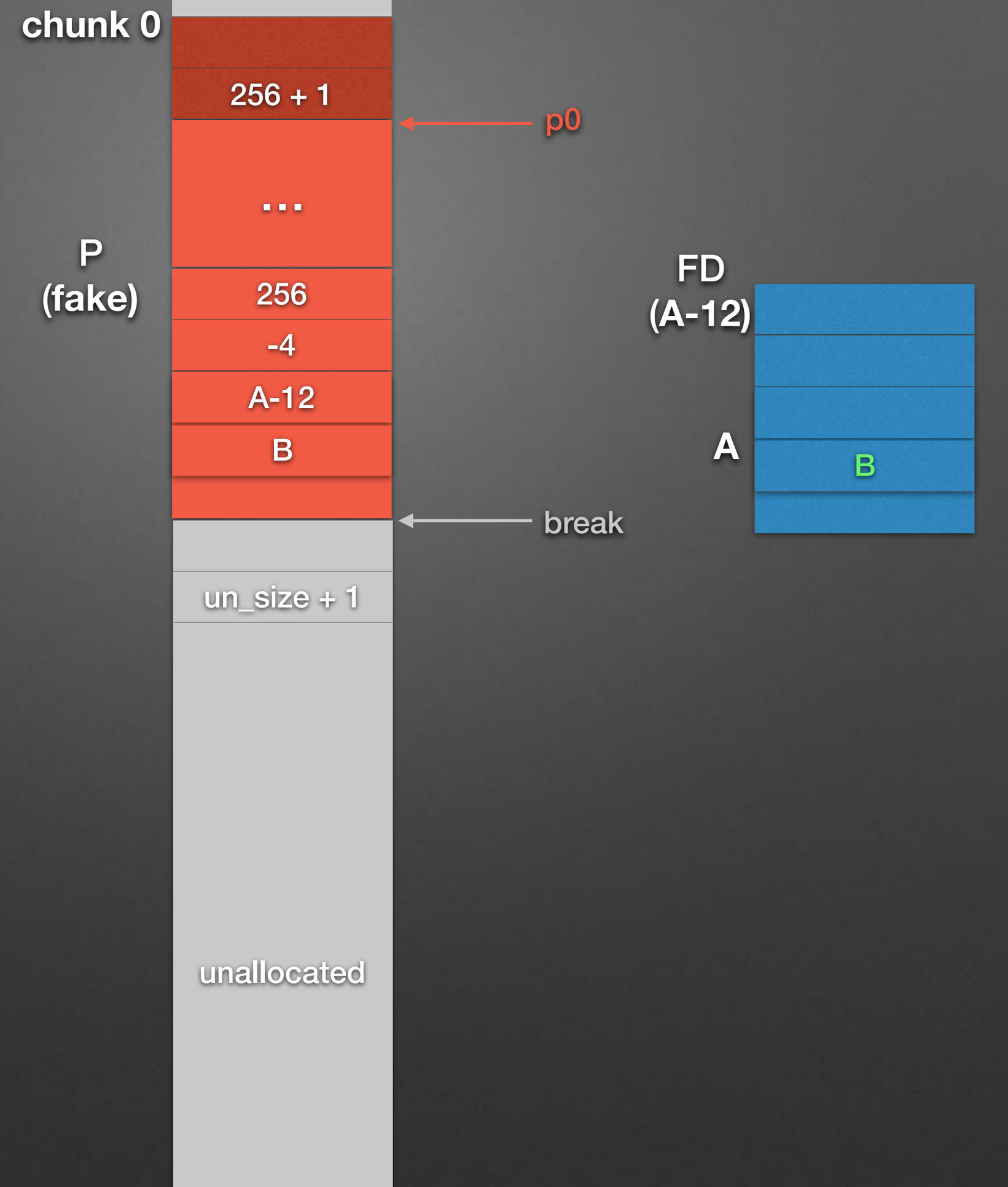
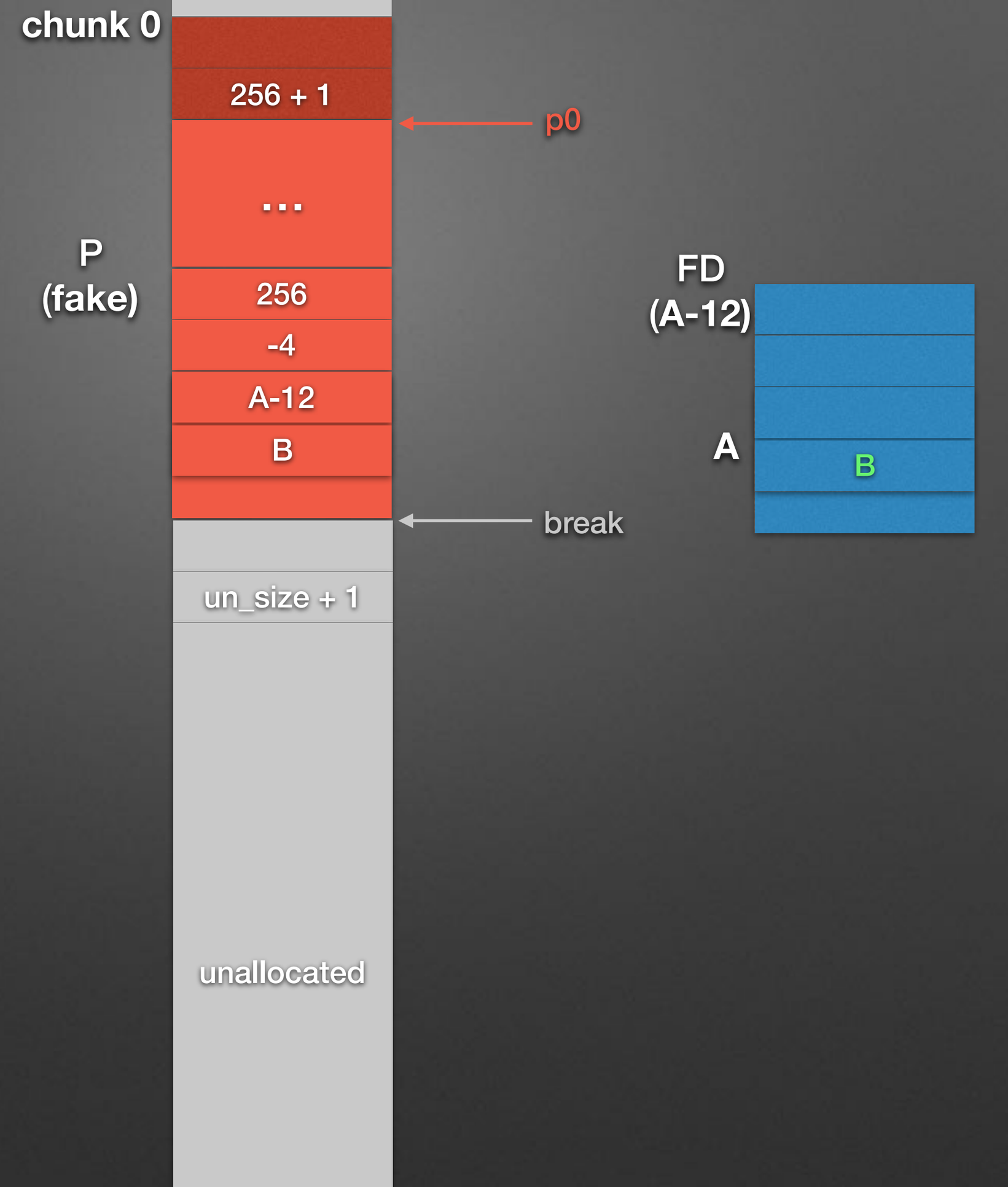

unlink(fake)




```
#define unlink(P) {
    FD = P->fd;
    BK = P->bk;
    FD->bk = BK;
    BK->fd = FD;
    ...
}
```

unlink(fake)



实现了在地址A写入B，
由于A和B可以为任意
值，即可以实现任意地
址任意写。
副作用是会在B+8位置
写入A-12。