```
// list.cc
//
//
      Routines to manage a singly-linked list of "things".
//
// A "ListElement" is allocated for each item to be put on the
//list; it is de-allocated when the item is removed. This means
      we don't need to keep a "next" pointer in every object we
//
      want to put on a list.
     NOTE: Mutual exclusion must be provided by the caller.
// If you want a synchronized list, you must use the routines
//in synchlist.cc.
#include "copyright.h"
#include "list.h"
// ListElement::ListElement
// Initialize a list element, so it can be added somewhere on a list.
//"itemPtr" is the item to be put on the list. It can be a pointer
//to anything.
//"sortKey" is the priority of the item, if any.
//-----
ListElement::ListElement(void *itemPtr, int sortKey)
    item = itemPtr;
    key = sortKey;
    next = NULL;// assume we'll put it at the end of the list
//-----
// List::List
//Initialize a list, empty to start with.
//Elements can now be added to the list.
//----
List::List()
   first = last = NULL;
´//-----
// List::~List
//Prepare a list for deallocation. If the list still contains any
//ListElements, de-allocate them. However, note that we do *not*
//de-allocate the "items" on the list -- this module allocates
//and de-allocates the ListElements to keep track of each item,
//but a given item may be on multiple lists, so we can't
//de-allocate them here.
List::~List()
   while (Remove() != NULL)
; // delete all the list elements
.//-----
// List::Append
      Append an "item" to the end of the list.
//Allocate a ListElement to keep track of the item.
      If the list is empty, then this will be the only element.
//Otherwise, put it at the end.
```

```
//"item" is the thing to put on the list, it can be a pointer to
anvthing.
//-----
List::Append(void *item)
   ListElement *element = new ListElement(item, 0);
   if (IsEmpty()) {// list is empty
first = element;
last = element;
   } else {// else put it after last
last->next = element;
last = element;
//----
// List::Prepend
// Put an "item" on the front of the list.
//Allocate a ListElement to keep track of the item.
      If the list is empty, then this will be the only element.
//Otherwise, put it at the beginning.
//"item" is the thing to put on the list, it can be a pointer to
//----
void
List::Prepend(void *item)
   ListElement *element = new ListElement(item, 0);
   if (IsEmpty()) {// list is empty
first = element;
last = element;
   } else {// else put it before first
element->next = first;
first = element;
    _____
// List::Remove
      Remove the first "item" from the front of the list.
// Returns:
//Pointer to removed item, NULL if nothing on the list.
void *
List::Remove()
   return SortedRemove(NULL); // Same as SortedRemove, but ignore
the key
//-----
// List::Mapcar
//Apply a function to each item on the list, by walking through
//the list, one element at a time.
//Unlike LISP, this mapcar does not return anything!
//"func" is the procedure to apply to each element of the list.
void
```

```
List::Mapcar(VoidFunctionPtr func)
   for (ListElement *ptr = first; ptr != NULL; ptr = ptr->next) {
      DEBUG('l', "In mapcar, about to invoke %x(%x)\n", func. ptr-
>item);
      (*func)((int)ptr->item);
       _____
// List::IsEmptv
      Returns TRUE if the list is empty (has no items).
bool
List::IsEmpty()
   if (first == NULL)
       return TRUE;
   else
return FALSE;
            .____
// List::SortedInsert
       Insert an "item" into a list, so that the list elements are
//sorted in increasing order by "sortKey".
//Allocate a ListElement to keep track of the item.
       If the list is empty, then this will be the only element.
//Otherwise, walk through the list, one element at a time,
//to find where the new item should be placed.
//"item" is the thing to put on the list, it can be a pointer to
//anything.
//"sortKey" is the priority of the item.
List::SortedInsert(void *item, int sortKey)
   ListElement *element = new ListElement(item, sortKey);
   ListElement *ptr;// keep track
   if (IsEmpty()) {// if list is empty, put
       first = element;
       last = element;
   } else if (sortKey < first->key) {
// item goes on front of list
element->next = first;
first = element;
   } else {// look for first elt in list bigger than item
       for (ptr = first; ptr->next != NULL; ptr = ptr->next) {
           if (sortKey < ptr->next->key) {
element->next = ptr->next;
       ptr->next = element;
return;
last->next = element;// item goes at end of list
last = element;
```

```
// List::SortedRemove
        Remove the first "item" from the front of a sorted list.
// Returns:
//Pointer to removed item, NULL if nothing on the list.
//Sets *keyPtr to the priority value of the removed item
//(this is needed by interrupt.cc, for instance).
//"keyPtr" is a pointer to the location in which to store the
//priority of the removed item.
void *
List::SortedRemove(int *keyPtr)
    ListElement *element = first;
    void *thing;
    if (IsEmpty())
return NULL;
    thing = first->item;
    if (first == last) {// list had one item, now has none
        first = NULL;
last = NULL;
    } else {
        first = element->next;
    if (keyPtr != NULL)
        *keyPtr = element->key;
    delete element;
    return thing;
```

```
// list.h
//Data structures to manage LISP-like lists.
//
       As in LISP, a list can contain any type of data structure
//as an item on the list: thread control blocks.
//pending interrupts, etc. That is why each item is a "void *",
//or in other words, a "pointers to anything".
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#ifndef LIST H
#define LIST_H
#include "copyright.h"
#include "utility.h"
// The following class defines a "list element" -- which is
// used to keep track of one item on a list. It is equivalent to a
// LISP cell, with a "car" ("next") pointing to the next element on the
list.
// and a "cdr" ("item") pointing to the item on the list.
// Internal data structures kept public so that List operations can
// access them directly.
class ListElement {
   public:
     ListElement(void *itemPtr, int sortKey);// initialize a list
element
     ListElement *next; // next element on list.
// NULL if this is the last
     int key;
                // priority, for a sorted list
     void *item;
                    // pointer to item on the list
// The following class defines a "list" -- a singly linked list of
// list elements, each of which points to a single item on the list.
// By using the "Sorted" functions, the list can be kept in sorted
// in increasing order by "key" in ListElement.
class List {
 public:
   List();// initialize the list
   ~List();// de-allocate the list
   void Prepend(void *item); // Put item at the beginning of the list
   void Append(void *item); // Put item at the end of the list
   void *Remove(); // Take item off the front of the list
   void Mapcar(VoidFunctionPtr func);// Apply "func" to every element
// on the list
   bool IsEmpty();// is the list empty?
   // Routines to put/get items on/off list in order (sorted by key)
   void SortedInsert(void *item, int sortKey);// Put item into list
   void *SortedRemove(int *keyPtr); // Remove first item from list
   ListElement *first; // Head of the list, NULL if list is empty
   ListElement *last; // Last element of list
```

#endif // LIST H

```
// main.cc
//Bootstrap code to initialize the operating system kernel.
//Allows direct calls into internal operating system functions,
//to simplify debugging and testing. In practice, the
//bootstrap code would just initialize data structures,
//and start a user program to print the login prompt.
//
// Most of this file is not needed until later assignments.
// Usage: nachos -d <debugflags> -rs <random seed #>
//-s -x <nachos file> -c <consoleIn> <consoleOut>
//-f -cp <unix file> <nachos file>
//-p <nachos file> -r <nachos file> -l -D -t
//
               -n <network reliability> -m <machine id>
//
               -o <other machine id>
//
               -z
//
//
      -d causes certain debugging messages to be printed (cf.
utility.h)
//
      -rs causes Yield to occur at random (but repeatable) spots
//
      -z prints the copyright message
//
   USER PROGRAM
//
      -s causes user programs to be executed in single-step mode
      -x runs a user program
//
      -c tests the console
//
   FILESYS
     -f causes the physical disk to be formatted
//
      -cp copies a file from UNIX to Nachos
//
      -p prints a Nachos file to stdout
//
      -r removes a Nachos file from the file system
      -1 lists the contents of the Nachos directory
//
//
      -D prints the contents of the entire file system
//
      -t tests the performance of the Nachos file system
//
// NETWORK
//
      -n sets the network reliability
      -m sets this machine's host id (needed for the network)
//
//
      -o runs a simple test of the Nachos network software
//
// NOTE -- flags are ignored until the relevant assignment.
   Some of the flags are interpreted here; some in system.cc.
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#define MAIN
#include "copyright.h"
#undef MAIN
#include "utility.h"
#include "system.h"
#ifdef THREADS
extern int testnum;
```

```
#endif
// External functions used by this file
extern void ThreadTest(void), Copy(char *unixFile, char *nachosFile);
extern void Print(char *file), PerformanceTest(void);
extern void StartProcess(char *file), ConsoleTest(char *in, char
extern void MailTest(int networkID);
// Bootstrap the operating system kernel.
//Check command line arguments
//Initialize data structures
//(optionally) Call test procedure
//"argc" is the number of command line arguments (including the name
//of the command) -- ex: "nachos -d +" -> argc = 3
//"argv" is an array of strings, one for each command line argument
//ex: "nachos -d +" -> argv = { "nachos", "-d", "+" }
int
main(int argc, char **argv)
    int argCount; // the number of arguments
// for a particular command
    DEBUG('t', "Entering main");
    (void) Initialize(argc, argv);
#ifdef THREADS
    for (argc--, argv++; argc > 0; argc -= argCount, argv += argCount)
      argCount = 1;
      switch (argv[0][1]) {
      case 'q':
        testnum = atoi(argv[1]);
        argCount++;
        break;
      default:
        testnum = 1;
        break;
    ThreadTest();
#endif
    for (argc--, argv++; argc > 0; argc -= argCount, argv += argCount)
argCount = 1;
        if (!strcmp(*argv, "-z"))
                                                // print copyright
            printf (copyright);
#ifdef USER PROGRAM
        if (!strcmp(*argv, "-x")) {
                                           // run a user program
    ASSERT(argc > 1);
            StartProcess(*(argv + 1));
            argCount = 2;
        } else if (!strcmp(*argv, "-c")) {
                                                // test the console
    if (argc == 1)
```

```
ConsoleTest(NULL, NULL);
    else {
ASSERT(argc > 2);
        ConsoleTest(*(argv + 1), *(argv + 2));
        argCount = 3;
    interrupt->Halt();// once we start the console, then
// Nachos will loop forever waiting
// for console input
#endif // USER_PROGRAM
#ifdef FILESYS
if (!strcmp(*argv, "-cp")) { // copy from UNIX to Nachos
   ASSERT(argc > 2);
    Copy(*(argv + 1), *(argv + 2));
    argCount = 3;
} else if (!strcmp(*argv, "-p")) {// print a Nachos file
   ASSERT(argc > 1);
    Print(*(argv + 1));
    argCount = 2;
} else if (!strcmp(*argv, "-r")) {// remove Nachos file
   ASSERT(argc > 1);
    fileSystem->Remove(*(argv + 1));
    argCount = 2;
} else if (!strcmp(*argv, "-1")) {// list Nachos directory
            fileSystem->List();
 else if (!strcmp(*argv, "-D")) {// print entire filesystem
            fileSystem->Print();
} else if (!strcmp(*argv, "-t")) {// performance test
            PerformanceTest();
#endif // FILESYS
#ifdef NETWORK
        if (!strcmp(*argv, "-o")) {
    ASSERT(argc > 1);
            Delay(2); // delay for 2 seconds
// to give the user time to
// start up another nachos
            MailTest(atoi(*(argv + 1)));
            argCount = 2;
#endif // NETWORK
    currentThread->Finish();// NOTE: if the procedure "main"
// returns, then the program "nachos"
// will exit (as any other normal program
// would). But there may be other
// threads on the ready list. We switch
// to those threads by saying that the
// "main" thread is finished, preventing
// it from returning.
    return(0);// Not reached...
```

```
// scheduler.cc
//Routines to choose the next thread to run, and to dispatch to
//that thread.
//
// These routines assume that interrupts are already disabled.
//If interrupts are disabled, we can assume mutual exclusion
//(since we are on a uniprocessor).
//
// NOTE: We can't use Locks to provide mutual exclusion here, since
// if we needed to wait for a lock, and the lock was busy, we would
//end up calling FindNextToRun(), and that would put us in an
//infinite loop.
//
// Very simple implementation -- no priorities, straight FIFO.
//Might need to be improved in later assignments.
//
// Copyright (c) 1992-1993 The Regents of the University of California. //
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "scheduler.h"
#include "system.h"
//-----
// Scheduler::Scheduler
// Initialize the list of ready but not running threads to empty.
//----
Scheduler::Scheduler()
   readyList = new List;
  ______
// Scheduler::~Scheduler
// De-allocate the list of ready threads.
Scheduler::~Scheduler()
   delete readyList;
// Scheduler::ReadyToRun
// Mark a thread as ready, but not running.
//Put it on the ready list, for later scheduling onto the CPU.
//"thread" is the thread to be put on the ready list.
//-----
void
Scheduler::ReadyToRun (Thread *thread)
   DEBUG('t', "Putting thread %s on ready list.\n", thread-
>qetName());
   thread->setStatus(READY);
   readyList->Append((void *)thread);
//-----
// Scheduler::FindNextToRun
```

```
// Return the next thread to be scheduled onto the CPU.
//If there are no ready threads, return NULL.
// Side effect:
//Thread is removed from the ready list.
Thread *
Scheduler::FindNextToRun ()
   return (Thread *)readyList->Remove();
//----
// Scheduler::Run
// Dispatch the CPU to nextThread. Save the state of the old thread,
//and load the state of the new thread, by calling the machine
//dependent context switch routine, SWITCH.
11
      Note: we assume the state of the previously running thread has
//already been changed from running to blocked or ready (depending).
// Side effect:
//The global variable currentThread becomes nextThread.
//"nextThread" is the thread to be put into the CPU.
void
Scheduler::Run (Thread *nextThread)
   Thread *oldThread = currentThread;
#ifdef USER PROGRAM// ignore until running user programs
   if (currentThread->space != NULL) {// if this thread is a user
program,
       currentThread->SaveUserState(); // save the user's CPU
registers
currentThread->space->SaveState();
#endif
   oldThread->CheckOverflow(); // check if the old thread
   // had an undetected stack overflow
   currentThread = nextThread;
                             // switch to the next thread
   running
   DEBUG('t', "Switching from thread \"%s\" to thread \"%s\"\n",
 oldThread->getName(), nextThread->getName());
   // This is a machine-dependent assembly language routine defined
   // in switch.s. You may have to think
   // a bit to figure out what happens after this, both from the
point
   // of view of the thread and from the perspective of the "outside
world".
   SWITCH(oldThread, nextThread);
   DEBUG('t', "Now in thread \"%s\"\n", currentThread->getName());
   // If the old thread gave up the processor because it was
```

```
finishing,
   // we need to delete its carcass. Note we cannot delete the thread
   // before now (for example, in Thread::Finish()), because up to
this
   // point, we were still running on the old thread's stack!
   if (threadToBeDestroyed != NULL) {
       delete threadToBeDestroyed;
threadToBeDestroyed = NULL;
#ifdef USER PROGRAM
   if (currentThread->space != NULL) {// if there is an address space
       currentThread->RestoreUserState();  // to restore, do it.
currentThread->space->RestoreState();
#endif
//-----
// Scheduler::Print
// Print the scheduler state -- in other words, the contents of
//the ready list. For debugging.
void
Scheduler::Print()
   printf("Ready list contents:\n");
   readyList->Mapcar((VoidFunctionPtr) ThreadPrint);
```

```
// scheduler.h
//Data structures for the thread dispatcher and scheduler.
//Primarily, the list of threads that are ready to run.
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#ifndef SCHEDULER H
#define SCHEDULER H
#include "copyright.h"
#include "list.h"
#include "thread.h"
// The following class defines the scheduler/dispatcher abstraction --
// the data structures and operations needed to keep track of which
// thread is running, and which threads are ready but not running.
class Scheduler {
 public:
   Scheduler();// Initialize list of ready threads
   ~Scheduler();// De-allocate ready list
   void ReadyToRun(Thread* thread);// Thread can be dispatched.
   Thread* FindNextToRun();// Dequeue first thread on the ready
// list, if any, and return thread.
   void Run(Thread* nextThread);// Cause nextThread to start running
   void Print();// Print contents of ready list
 private:
   List *readyList; // queue of threads that are ready to run,
// but not running
};
#endif // SCHEDULER H
```

```
/* switch.s
     Machine dependent context switch routines. DO NOT MODIFY THESE!
 *Context switching is inherently machine dependent, since
 *the registers to be saved, how to set up an initial
 *call frame, etc, are all specific to a processor architecture.
 * This file currently supports the following architectures:
      DEC MIPS
      SUN SPARC
      HP PA-RISC
      Intel 386
 * We define two routines for each architecture:
 * ThreadRoot(InitialPC, InitialArg, WhenDonePC, StartupPC)
 *InitialPC - The program counter of the procedure to run
 *in this thread.
        InitialArg - The single argument to the thread.
 *WhenDonePC - The routine to call when the thread returns.
 *StartupPC - Routine to call when the thread is started.
 *ThreadRoot is called from the SWITCH() routine to start
 *a thread for the first time.
 * SWITCH(oldThread, newThread)
 * oldThread - The current thread that was running, where the
 *CPU register state is to be saved.
 * newThread - The new thread to be run, where the CPU register
 *state is to be loaded from.
#include "copyright.h"
#include "switch.h"
#ifdef HOST MIPS
/* Symbolic register names */
#define z
                $0
                        /* zero register */
                        /* argument registers */
#define a0
                $4
#define al
                $5
#define s0
                $16
                        /* callee saved */
#define s1
                $17
#define s2
                $18
#define s3
                $19
                $20
#define s4
                $21
#define s5
#define s6
                $22
#define s7
                $23
#define sp
                $29
                        /* stack pointer */
#define fp
                $30
                        /* frame pointer */
                        /* return address */
#define ra
                $31
        .text
        .aliqn 2
.globl ThreadRoot
.entThreadRoot,0
ThreadRoot:
orfp,z,z# Clearing the frame pointer here
# makes qdb backtraces of thread stacks
```

```
# end here (I hope!)
ialStartupPC# call startup procedure
movea0, InitialArg
jalInitialPC# call main procedure
ial WhenDonePC# when were done, call clean up procedure
# NEVER REACHED
end ThreadRoot.
# a0 -- pointer to old Thread
# al -- pointer to new Thread
.qlobl SWITCH
.entSWITCH,0
SWITCH:
swsp, SP(a0)# save new stack pointer
sws0, S0(a0)# save all the callee-save registers
sws1, S1(a0)
sws2, S2(a0)
sws3, S3(a0)
sws4, S4(a0)
sws5, S5(a0)
sws6, S6(a0)
sws7, S7(a0)
swfp, FP(a0)# save frame pointer
swra, PC(a0)# save return address
lwsp, SP(a1)# load the new stack pointer
lws0, S0(a1) # load the callee-save registers
lws1, S1(a1)
lws2, S2(a1)
lws3, S3(a1)
lws4, S4(a1)
lws5, S5(a1)
lws6, S6(a1)
lws7, S7(a1)
lwfp, FP(a1)
lwra, PC(a1)# load the return address
ira
.end SWITCH
#endif // HOST_MIPS
#ifdef HOST SPARC
/* NOTE! deleted ......
 * These files appear not to exist on Solaris --
 * you need to find where (the SPARC-specific) MINFRAME,
ST FLUSH WINDOWS, ...
 * are defined. (I don't have a Solaris machine, so I have no way to
tell.)
 * /
/* From sys/trap.h and sys/asm linkage.h */
#endif // HOST_SPARC
#ifdef HOST SNAKE
/* NOTE! deleted ......
 * These files appear not to exist on Solaris --
 * you need to find where (the SPARC-specific) MINFRAME,
ST FLUSH WINDOWS, ...
 * are defined. (I don't have a Solaris machine, so I have no way to
tell.)
/* From sys/trap.h and sys/asm_linkage.h */
```

```
#endif
#ifdef HOST i386
        .text
        .align 2
        .globl ThreadRoot
   void ThreadRoot( void )
   expects the following registers to be initialized:
* *
                points to startup function (interrupt enable)
* *
        edx
                contains inital argument to thread function
**
                points to thread function
        esi
* *
        edi
                point to Thread::Finish()
* /
ThreadRoot:
        pushl
                %ebp
                %esp,%ebp
        movl
        pushl
                InitialArg
        call
                 *StartupPC
        call
                 *InitialPC
        call
                 *WhenDonePC
        // NOT REACHED
        movl
                %ebp,%esp
        popl
                %ebp
        ret
   void SWITCH( thread *t1, thread *t2 )
   on entry, stack looks like this:
* *
        8(esp)
                                 thread *t2
               ->
* *
        4(esp) ->
                                 thread *t1
                                 return address
         (esp) ->
* *
   we push the current eax on the stack so that we can use it as
   a pointer to t1, this decrements esp by 4, so when we use it
** to reference stuff on the stack, we add 4 to the offset.
        .comm
                _eax_save,4
        .globl SWITCH
SWITCH:
                %eax, eax save
                                          # save the value of eax
        movl
        movl
                4(%esp),%eax
                                         # move pointer to t1 into eax
        movl
                %ebx, EBX(%eax)
                                          # save registers
        movl
                %ecx,_ECX(%eax)
        movl
                %edx,_EDX(%eax)
        movl
                %esi,_ESI(%eax)
        movl
                %edi, EDI(%eax)
        movl
                %ebp,_EBP(%eax)
        movl
                %esp, ESP(%eax)
                                          # save stack pointer
                                          # get the saved value of eax
        movl
                 _eax_save, %ebx
        movl
                %ebx,_EAX(%eax)
                                          # store it
        movl
                0(%esp),%ebx
                                          # get return address from stack
into ebx
        movl
                %ebx,_PC(%eax)
                                          # save it into the pc storage
        movl
                8(%esp),%eax
                                          # move pointer to t2 into eax
        movl
                 _EAX(%eax),%ebx
                                          # get new value for eax into
ebx
```

```
2/2
                %ebx, eax save
                                        # save it
        movl
                EBX(%eax),%ebx
                                        # retore old registers
        movl
       movl
                ECX(%eax),%ecx
       movl
                EDX(%eax),%edx
       movl
                ESI(%eax),%esi
        movl
                EDI(%eax),%edi
                EBP(%eax),%ebp
       movl
                                         # restore stack pointer
        movl
                ESP(%eax),%esp
                                        # restore return address into
       movl
                _PC(%eax),%eax
eax
        movl
                %eax,4(%esp)
                                         # copy over the ret address on
the stack
        movl
                eax save, %eax
       ret
#endif
```

```
/* switch.h
                                                                         /* Registers that must be saved during a context switch. See comment
 *Definitions needed for implementing context switching.
                                                                         above. */
                                                                         #define I0 4
                                                                         #define I1 8
 *Context switching is inherently machine dependent, since
 *the registers to be saved, how to set up an initial
                                                                         #define T2 12
 *call frame, etc, are all specific to a processor architecture.
                                                                         #define I3 16
                                                                         #define I4 20
 * This file currently supports the DEC MIPS and SUN SPARC
                                                                         #define I5 24
architectures.
                                                                         #define I6 28
 * /
                                                                         #define I7 32
/*
                                                                         /* Aliases used for clearing code. */
 Copyright (c) 1992-1993 The Regents of the University of California.
                                                                         #define FP I6
All rights reserved. See copyright.h for copyright notice and
                                                                         #define PC I7
limitation
                                                                         /* Registers for ThreadRoot. See comment above. */
 of liability and disclaimer of warranty provisions.
                                                                         #define InitialPC
                                                                                                 %00
 * /
                                                                         #define InitialArg
                                                                                                 %o1
#ifndef SWITCH_H
                                                                         #define WhenDonePC
                                                                                                 %02
#define SWITCH_H
                                                                         #define StartupPC
                                                                                                 %03
#include "copyright.h"
                                                                         #define PCState
                                                                                                 (PC/4-1)
#ifdef HOST_MIPS
                                                                         #define InitialPCState (I0/4-1)
/* Registers that must be saved during a context switch.
                                                                         #define InitialArgState (I1/4-1)
 * These are the offsets from the beginning of the Thread object,
                                                                         #define WhenDonePCState (I2/4-1)
 * in bytes, used in switch.s
                                                                         #define StartupPCState (I3/4-1)
 * /
                                                                         #endif // HOST_SPARC
#define SP 0
                                                                         #ifdef HOST SNAKE
#define S0 4
                                                                         /* Registers that must be saved during a context switch. See comment
                                                                         above. */
#define S1 8
                                                                         #define
                                                                                   SP
#define S2 12
                                                                                        4
#define S3 16
                                                                         #define
                                                                                   S0
#define S4 20
                                                                         #define
                                                                                   S1
                                                                                        8
#define S5 24
                                                                                        12
                                                                         #define
                                                                                   S2
#define S6 28
                                                                         #define
                                                                                        16
                                                                                        20
#define S7 32
                                                                         #define
                                                                                   S4
#define FP 36
                                                                                        24
                                                                         #define
                                                                                   S_5
#define PC 40
                                                                         #define
                                                                                   S6
                                                                                        28
/* To fork a thread, we set up its saved register state, so that
                                                                         #define
                                                                                   S7
                                                                                        32
 * when we switch to the thread, it will start running in ThreadRoot.
                                                                         #define
                                                                                   S8
                                                                                        36
                                                                         #define
                                                                                   S9
                                                                                        40
 * The following are the initial registers we need to set up to
                                                                         #define
                                                                                   S10
                                                                                        44
 * pass values into ThreadRoot (for instance, containing the procedure |#define
                                                                                        48
                                                                                   S11
 * for the thread to run). The first set is the registers as used
                                                                         #define
                                                                                   S12 52
 * by ThreadRoot; the second set is the locations for these initial
                                                                         #define
                                                                                   S13
                                                                                        56
 * values in the Thread object -- used in Thread::AllocateStack().
                                                                         #define
                                                                                   S14 60
                                                                         #define
                                                                                   S15 64
#define InitialPCs0
                                                                         #define
                                                                                   PC
                                                                                        68
                                                                         /* Registers for ThreadRoot.
#define InitialArgs1
                                                                                                       See comment above. */
#define WhenDonePCs2
                                                                         #define InitialPC
                                                                                                 %r3/* S0 */
#define StartupPCs3
                                                                         #define InitialArg
                                                                                                 %r4
#define PCState(PC/4-1)
                                                                         #define WhenDonePC
                                                                                                 %r5
#define FPState(FP/4-1)
                                                                         #define StartupPC
                                                                                                 %r6
#define InitialPCState(S0/4-1)
                                                                         #define PCState
                                                                                                  (PC/4-1)
                                                                         #define InitialPCState (S0/4-1)
#define InitialArgState(S1/4-1)
#define WhenDonePCState(S2/4-1)
                                                                         #define InitialArgState (S1/4-1)
#define StartupPCState(S3/4-1)
                                                                         #define WhenDonePCState (S2/4-1)
#endif // HOST_MIPS
                                                                         #define StartupPCState (S3/4-1)
#ifdef HOST_SPARC
                                                                         #endif // HOST_SNAKE
```

```
#ifdef HOST i386
/* the offsets of the registers from the beginning of the thread object
#define _ESP
#define EAX
                  4
#define _EBX
                  8
#define _ECX
                 12
#define _EDX
                 16
                 20
#define _EBP
                  24
#define _ESI
#define _EDI
                  28
                 32
#define _PC
/* These definitions are used in Thread::AllocateStack(). */
#define PCState
                         (PC/4-1)
#define FPState
                         (_{EBP}/4-1)
#define InitialPCState
                         (_{ESI}/4-1)
#define InitialArgState (_EDX/4-1)
#define WhenDonePCState (_EDI/4-1)
#define StartupPCState
                         (\underline{ECX}/4-1)
#define InitialPC
                         %esi
#define InitialArg
                         %edx
#define WhenDonePC
                         %edi
#define StartupPC
                         %ecx
#endif
#endif // SWITCH_H
```

```
// synch.cc
//Routines for synchronizing threads. Three kinds of
//synchronization routines are defined here: semaphores, locks
    and condition variables (the implementation of the last two
//are left to the reader).
// Any implementation of a synchronization routine needs some
// primitive atomic operation. We assume Nachos is running on
// a uniprocessor, and thus atomicity can be provided by
// turning off interrupts. While interrupts are disabled, no
// context switch can occur, and thus the current thread is guaranteed
// to hold the CPU throughout, until interrupts are reenabled.
// Because some of these routines might be called with interrupts
// already disabled (Semaphore:: V for one), instead of turning
// on interrupts at the end of the atomic operation, we always simply
// re-set the interrupt state back to its original value (whether
// that be disabled or enabled).
#include "copyright.h"
#include "synch.h"
#include "system.h"
// Semaphore::Semaphore
// Initialize a semaphore, so that it can be used for synchronization.
//"debugName" is an arbitrary name, useful for debugging.
//"initialValue" is the initial value of the semaphore.
//-----
Semaphore::Semaphore(char* debugName, int initialValue)
   name = debugName;
   value = initialValue;
   queue = new List;
.//-----
// Semaphore::Semaphore
// De-allocate semaphore, when no longer needed. Assume no one
//is still waiting on the semaphore!
//-----
Semaphore::~Semaphore()
   delete queue;
.
//-----
// Semaphore::P
// Wait until semaphore value > 0, then decrement. Checking the
//value and decrementing must be done atomically, so we
//need to disable interrupts before checking the value.
//Note that Thread::Sleep assumes that interrupts are disabled
//when it is called.
void
Semaphore::P()
   IntStatus oldLevel = interrupt->SetLevel(IntOff);// disable
interrupts
   while (value == 0) { // semaphore not available
```

```
queue->Append((void *)currentThread);// so go to sleep
currentThread->Sleep();
    value--; // semaphore available,
// consume its value
    (void) interrupt->SetLevel(oldLevel);// re-enable interrupts
// Semaphore::V
// Increment semaphore value, waking up a waiter if necessary.
//As with P(), this operation must be atomic, so we need to disable
//interrupts. Scheduler::ReadyToRun() assumes that threads
//are disabled when it is called.
Semaphore::V()
    Thread *thread;
    IntStatus oldLevel = interrupt->SetLevel(IntOff);
    thread = (Thread *)queue->Remove();
    if (thread != NULL) // make thread ready, consuming the V
immediately
scheduler->ReadyToRun(thread);
    value++;
    (void) interrupt->SetLevel(oldLevel);
// Dummy functions -- so we can compile our later assignments
// Note -- without a correct implementation of Condition::Wait(),
// the test case in the network assignment won't work!
Lock::Lock(char* debugName) {}
Lock::~Lock() {}
void Lock::Acquire() {}
void Lock::Release() {}
Condition::Condition(char* debugName) { }
Condition::~Condition() { }
void Condition::Wait(Lock* conditionLock) { ASSERT(FALSE); }
void Condition::Signal(Lock* conditionLock) { }
void Condition::Broadcast(Lock* conditionLock) { }
```

```
// synch.h
//Data structures for synchronizing threads.
//Three kinds of synchronization are defined here: semaphores,
//locks, and condition variables. The implementation for
//semaphores is given; for the latter two, only the procedure
//interface is given -- they are to be implemented as part of
//the first assignment.
//Note that all the synchronization objects take a "name" as
//part of the initialization. This is solely for debugging purposes.
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// synch.h -- synchronization primitives.
#ifndef SYNCH H
#define SYNCH H
#include "copyright.h"
#include "thread.h"
#include "list.h"
// The following class defines a "semaphore" whose value is a non-
// integer. The semaphore has only two operations P() and V():
//P() -- waits until value > 0, then decrement
//
//V() -- increment, waking up a thread waiting in P() if necessary
// Note that the interface does *not* allow a thread to read the value
// the semaphore directly -- even if you did read the value, the
// only thing you would know is what the value used to be. You don't
// know what the value is now, because by the time you get the value
// into a register, a context switch might have occurred,
// and some other thread might have called P or V, so the true value
might
// now be different.
class Semaphore {
 public:
   Semaphore(char* debuqName, int initialValue);// set initial value
   ~Semaphore(); // de-allocate semaphore
   char* getName() { return name;}// debugging assist
   void P(); // these are the only operations on a semaphore
   void V(); // they are both *atomic*
 private:
   char* name;
                       // useful for debugging
                       // semaphore value, always >= 0
   int value;
   List *queue;
                       // threads waiting in P() for the value to be >
0
};
// The following class defines a "lock". A lock can be BUSY or FREE.
// There are only two operations allowed on a lock:
```

```
//Acquire -- wait until the lock is FREE, then set it to BUSY
//
//Release -- set lock to be FREE, waking up a thread waiting
//in Acquire if necessary
11
// In addition, by convention, only the thread that acquired the lock
// may release it. As with semaphores, you can't read the lock value
// (because the value might change immediately after you read it).
class Lock {
  public:
    Lock(char* debugName); // initialize lock to be FREE
    ~Lock();// deallocate lock
    char* getName() { return name; }// debugging assist
    void Acquire(); // these are the only operations on a lock
    void Release(); // they are both *atomic*
    bool isHeldByCurrentThread();// true if the current thread
// holds this lock. Useful for
// checking in Release, and in
// Condition variable ops below.
  private:
    char* name; // for debugging
    // plus some other stuff you'll need to define
// The following class defines a "condition variable". A condition
// variable does not have a value, but threads may be queued, waiting
// on the variable. These are only operations on a condition
variable:
//Wait() -- release the lock, relinquish the CPU until signaled,
//then re-acquire the lock
//Signal() -- wake up a thread, if there are any waiting on
//the condition
//Broadcast() -- wake up all threads waiting on the condition
// All operations on a condition variable must be made while
// the current thread has acquired a lock. Indeed, all accesses
// to a given condition variable must be protected by the same lock.
// In other words, mutual exclusion must be enforced among threads
calling
// the condition variable operations.
//
// In Nachos, condition variables are assumed to obey *Mesa*-style
// semantics. When a Signal or Broadcast wakes up another thread,
// it simply puts the thread on the ready list, and it is the
responsibility
// of the woken thread to re-acquire the lock (this re-acquire is
// taken care of within Wait()). By contrast, some define condition
// variables according to *Hoare*-style semantics -- where the
signalling
// thread gives up control over the lock and the CPU to the woken
// which runs immediately and gives back control over the lock to the
// signaller when the woken thread leaves the critical section.
```

```
// The consequence of using Mesa-style semantics is that some other
thread
// can acquire the lock, and change data structures, before the woken
// thread gets a chance to run.
class Condition {
 public:
   Condition(char* debugName);// initialize condition to
// "no one waiting"
   ~Condition();// deallocate the condition
   char* getName() { return (name); }
   void Wait(Lock *conditionLock); // these are the 3 operations on
// condition variables; releasing the
// lock and going to sleep are
// *atomic* in Wait()
   void Signal(Lock *conditionLock);  // conditionLock must be held
   void Broadcast(Lock *conditionLock);// the currentThread for all of
// these operations
 private:
   char* name;
   // plus some other stuff you'll need to define
#endif // SYNCH H
```

```
// synchlist.cc
//Routines for synchronized access to a list.
//Implemented by surrounding the List abstraction
//with synchronization routines.
// Implemented in "monitor"-style -- surround each procedure with a
// lock acquire and release pair, using condition signal and wait for
// synchronization.
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "synchlist.h"
//-----
// SvnchList::SvnchList
//Allocate and initialize the data structures needed for a
//synchronized list, empty to start with.
//Elements can now be added to the list.
SynchList::SynchList()
   list = new List();
   lock = new Lock("list lock");
   listEmpty = new Condition("list empty cond");
.
.
.
// SynchList::~SynchList
//De-allocate the data structures created for synchronizing a list.
//-----
SynchList::~SynchList()
   delete list;
   delete lock;
   delete listEmpty;
// SynchList::Append
      Append an "item" to the end of the list. Wake up anyone
//waiting for an element to be appended.
//"item" is the thing to put on the list, it can be a pointer to
//------
void
SynchList::Append(void *item)
   lock->Acquire();// enforce mutual exclusive access to the list
   list->Append(item);
   listEmpty->Signal(lock);// wake up a waiter, if any
   lock->Release();
                -----
// SynchList::Remove
```

```
Remove an "item" from the beginning of the list. Wait if
//the list is empty.
// Returns:
//The removed item.
* biov
SynchList::Remove()
   void *item;
   lock->Acquire();// enforce mutual exclusion
   while (list->IsEmpty())
listEmpty->Wait(lock);// wait until list isn't empty
   item = list->Remove();
   ASSERT(item != NULL);
   lock->Release();
   return item;
     -----
// SynchList::Mapcar
       Apply function to every item on the list. Obey mutual
exclusion
//constraints.
//"func" is the procedure to be applied.
biov
SynchList::Mapcar(VoidFunctionPtr func)
   lock->Acquire();
   list->Mapcar(func);
   lock->Release();
```

```
// synchlist.h
//Data structures for synchronized access to a list.
//Implemented by surrounding the List abstraction
//with synchronization routines.
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
// of liability and disclaimer of warranty provisions.
#ifndef SYNCHLIST H
#define SYNCHLIST H
#include "copyright.h"
#include "list.h"
#include "synch.h"
// The following class defines a "synchronized list" -- a list for
which:
// these constraints hold:
//1. Threads trying to remove an item from a list will
//wait until the list has an element on it.
//2. One thread at a time can access list data structures
class SynchList {
 public:
   SynchList();// initialize a synchronized list
   ~SynchList();// de-allocate a synchronized list
   void Append(void *item);// append item to the end of the list,
// and wake up any thread waiting in remove
   void *Remove();// remove the first item from the front of
// the list, waiting if the list is empty
// apply function to every item in the list
   void Mapcar(VoidFunctionPtr func);
 private:
   List *list;// the unsynchronized list
   Lock *lock;// enforce mutual exclusive access to the list
   Condition *listEmpty;// wait in Remove if the list is empty
#endif // SYNCHLIST_H
```

```
// system.cc
//Nachos initialization and cleanup routines.
// Copyright (c) 1992-1993 The Regents of the University of California. //
// All rights reserved. See copyright.h for copyright notice and
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "system.h"
// This defines *all* of the global data structures used by Nachos.
// These are all initialized and de-allocated by this file.
Thread *currentThread;// the thread we are running now
Thread *threadToBeDestroyed; // the thread that just finished
Scheduler *scheduler;// the ready list
Interrupt *interrupt;// interrupt status
Statistics *stats;// performance metrics
Timer *timer; // the hardware timer device.
// for invoking context switches
#ifdef FILESYS NEEDED
FileSystem *fileSystem;
#endif
#ifdef FILESYS
SynchDisk *synchDisk;
#endif
#ifdef USER PROGRAM// requires either FILESYS or FILESYS STUB
Machine *machine; // user program memory and registers
#endif
#ifdef NETWORK
PostOffice *postOffice;
#endif
// External definition, to allow us to take a pointer to this function
extern void Cleanup();
//----
// TimerInterruptHandler
// Interrupt handler for the timer device. The timer device is
//set up to interrupt the CPU periodically (once every TimerTicks).
//This routine is called each time there is a timer interrupt,
//with interrupts disabled.
//
//Note that instead of calling Yield() directly (which would
//suspend the interrupt handler, not the interrupted thread
//which is what we wanted to context switch), we set a flag
//so that once the interrupt handler is done, it will appear as
//if the interrupted thread called Yield at the point it is
//was interrupted.
//
//"dummy" is because every interrupt handler takes one argument,
//whether it needs it or not.
static void
TimerInterruptHandler(int dummy)
   if (interrupt->getStatus() != IdleMode)
interrupt->YieldOnReturn();
```

```
// Initialize
// Initialize Nachos global data structures. Interpret command
//line arguments in order to determine flags for the initialization.
//"argc" is the number of command line arguments (including the name
//of the command) -- ex: "nachos -d +" -> argc = 3
//"argy" is an array of strings, one for each command line argument
//ex: "nachos -d +" -> argv = { "nachos", "-d", "+" }
void
Initialize(int argc, char **argv)
    int argCount;
    char* debugArgs = "";
    bool randomYield = FALSE;
#ifdef USER PROGRAM
    bool debugUserProg = FALSE;// single step user program
#endif
#ifdef FILESYS NEEDED
    bool format = FALSE;// format disk
#endif
#ifdef NETWORK
    double rely = 1;// network reliability
    int netname = 0;// UNIX socket name
#endif
    for (argc--, argv++; argc > 0; argc -= argCount, argv += argCount)
argCount = 1;
if (!strcmp(*argv, "-d")) {
    if (argc == 1)
debugArgs = "+";// turn on all debug flags
    debugArgs = *(argv + 1);
    argCount = 2;
} else if (!strcmp(*argv, "-rs")) {
    ASSERT(argc > 1);
    RandomInit(atoi(*(argv + 1)));// initialize pseudo-random
// number generator
    randomYield = TRUE;
    argCount = 2;
#ifdef USER_PROGRAM
if (!strcmp(*argv, "-s"))
    debugUserProg = TRUE;
#endif
#ifdef FILESYS NEEDED
if (!strcmp(*argv, "-f"))
    format = TRUE;
#endif
#ifdef NETWORK
if (!strcmp(*argv, "-1")) {
    ASSERT(argc > 1);
    rely = atof(*(argv + 1));
    argCount = 2;
```

```
else if (!strcmp(*arqv, "-m")) {
    ASSERT(argc > 1);
    netname = atoi(*(argv + 1));
    argCount = 2;
#endif
    DebugInit(debugArgs);// initialize DEBUG messages
    stats = new Statistics();// collect statistics
    interrupt = new Interrupt;// start up interrupt handling
    scheduler = new Scheduler();// initialize the ready queue
    if (randomYield)// start the timer (if needed)
timer = new Timer(TimerInterruptHandler, 0, randomYield);
    threadToBeDestroyed = NULL;
    // We didn't explicitly allocate the current thread we are running
in.
    // But if it ever tries to give up the CPU, we better have a Thread
    // object to save its state.
    currentThread = new Thread("main");
    currentThread->setStatus(RUNNING);
    interrupt->Enable();
    CallOnUserAbort(Cleanup);// if user hits ctl-C
#ifdef USER_PROGRAM
    machine = new Machine(debugUserProg);// this must come first
#endif
#ifdef FILESYS
    synchDisk = new SynchDisk("DISK");
#endif
#ifdef FILESYS NEEDED
    fileSystem = new FileSystem(format);
#endif
#ifdef NETWORK
    postOffice = new PostOffice(netname, rely, 10);
#endif
// Cleanup
// Nachos is halting. De-allocate global data structures.
void
Cleanup()
    printf("\nCleaning up...\n");
#ifdef NETWORK
    delete postOffice;
#endif
#ifdef USER PROGRAM
    delete machine;
#endif
#ifdef FILESYS NEEDED
    delete fileSystem;
#endif
#ifdef FILESYS
    delete synchDisk;
```

#endif delete timer; delete scheduler; delete interrupt; Exit(0);

```
// system.h
//All global variables used in Nachos are defined here.
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
// of liability and disclaimer of warranty provisions.
#ifndef SYSTEM H
#define SYSTEM_H
#include "copyright.h"
#include "utility.h"
#include "thread.h"
#include "scheduler.h"
#include "interrupt.h"
#include "stats.h"
#include "timer.h"
// Initialization and cleanup routines
extern void Initialize(int argc, char **argv); // Initialization,
// called before anything else
extern void Cleanup();// Cleanup, called when
// Nachos is done.
extern Thread *currentThread;// the thread holding the CPU
extern Thread *threadToBeDestroyed; // the thread that just finished
extern Scheduler *scheduler;// the ready list
extern Interrupt *interrupt;// interrupt status
extern Statistics *stats;// performance metrics
extern Timer *timer; // the hardware alarm clock
#ifdef USER PROGRAM
#include "machine.h"
extern Machine* machine; // user program memory and registers
#ifdef FILESYS NEEDED // FILESYS or FILESYS STUB
#include "filesys.h"
extern FileSystem *fileSystem;
#endif
#ifdef FILESYS
#include "synchdisk.h"
extern SynchDisk
                  *synchDisk;
#endif
#ifdef NETWORK
#include "post.h"
extern PostOffice* postOffice;
#endif
#endif // SYSTEM H
```

```
// thread.cc
//Routines to manage threads. There are four main operations:
//Fork -- create a thread to run a procedure concurrently
//with the caller (this is done in two steps -- first
//allocate the Thread object, then call Fork on it)
//Finish -- called when the forked procedure finishes, to clean up
//Yield -- relinquish control over the CPU to another ready thread
//Sleep -- relinguish control over the CPU, but thread is now blocked.
//In other words, it will not run again, until explicitly
//put back on the ready queue.
#include "copyright.h"
#include "thread.h"
#include "switch.h"
#include "synch.h"
#include "system.h"
#define STACK_FENCEPOST 0xdeadbeef// this is put at the top of the
// execution stack, for detecting stack overflows
// Thread::Thread
// Initialize a thread control block, so that we can then call
//Thread::Fork.
//"threadName" is an arbitrary string, useful for debugging.
Thread::Thread(char* threadName)
   name = threadName;
   stackTop = NULL;
   stack = NULL;
   status = JUST CREATED;
#ifdef USER PROGRAM
   space = NULL;
#endif
// Thread::~Thread
// De-allocate a thread.
// NOTE: the current thread *cannot* delete itself directly,
//since it is still running on the stack that we need to delete.
//
       NOTE: if this is the main thread, we can't delete the stack
//
       because we didn't allocate it -- we got it automatically
       as part of starting up Nachos.
Thread::~Thread()
   DEBUG('t', "Deleting thread \"%s\"\n", name);
   ASSERT(this != currentThread);
   if (stack != NULL)
DeallocBoundedArray((char *) stack, StackSize * sizeof(int));
// Thread::Fork
// Invoke (*func)(arg), allowing caller and callee to execute
//NOTE: although our definition allows only a single integer argument
//to be passed to the procedure, it is possible to pass multiple
//arguments by making them fields of a structure, and passing a pointer //or the execution stack, because we're still running in the thread
```

```
//to the structure as "arq".
11
// Implemented as the following steps:
//1. Allocate a stack
//2. Initialize the stack so that a call to SWITCH will
//cause it to run the procedure
//3. Put the thread on the ready queue
//
//"func" is the procedure to run concurrently.
//"arg" is a single argument to be passed to the procedure.
//----
Thread::Fork(VoidFunctionPtr func, int arg)
   DEBUG('t', "Forking thread \"%s\" with func = 0x%x, arg = %d\n",
 name, (int) func, arg);
   StackAllocate(func, arg);
   IntStatus oldLevel = interrupt->SetLevel(IntOff);
   scheduler->ReadyToRun(this);// ReadyToRun assumes that interrupts
// are disabled!
   (void) interrupt->SetLevel(oldLevel);
// Thread::CheckOverflow
// Check a thread's stack to see if it has overrun the space
//that has been allocated for it. If we had a smarter compiler,
//we wouldn't need to worry about this, but we don't.
//
// NOTE: Nachos will not catch all stack overflow conditions.
//In other words, your program may still crash because of an overflow.
// If you get bizarre results (such as seg faults where there is no
code)
// then you *may* need to increase the stack size. You can avoid
stack
// overflows by not putting large data structures on the stack.
// Don't do this: void foo() { int bigArray[10000]; ... }
//-----
void
Thread::CheckOverflow()
   if (stack != NULL)
#ifdef HOST SNAKE// Stacks grow upward on the Snakes
ASSERT(stack[StackSize - 1] == STACK FENCEPOST);
#else
ASSERT((int) *stack == (int) STACK FENCEPOST);
#endif
// Thread::Finish
// Called by ThreadRoot when a thread is done executing the
//forked procedure.
//
// NOTE: we don't immediately de-allocate the thread data structure
```

```
//and we're still on the stack! Instead, we set "threadToBeDestroyed",
//so that Scheduler::Run() will call the destructor, once we're
//running in the context of a different thread.
//
// NOTE: we disable interrupts, so that we don't get a time slice
//between setting threadToBeDestroyed, and going to sleep.
biov
Thread::Finish ()
   (void) interrupt->SetLevel(IntOff);
   ASSERT(this == currentThread);
   DEBUG('t', "Finishing thread \"%s\"\n", getName());
   threadToBeDestroyed = currentThread;
   Sleep();// invokes SWITCH
   // not reached
// Thread::Yield
// Relinguish the CPU if any other thread is ready to run.
//If so, put the thread on the end of the ready list, so that
//it will eventually be re-scheduled.
//NOTE: returns immediately if no other thread on the ready queue.
//Otherwise returns when the thread eventually works its way
//to the front of the ready list and gets re-scheduled.
//NOTE: we disable interrupts, so that looking at the thread
//on the front of the ready list, and switching to it, can be done
//atomically. On return, we re-set the interrupt level to its
//original state, in case we are called with interrupts disabled.
// Similar to Thread::Sleep(), but a little different.
//-----
void
Thread::Yield ()
   Thread *nextThread;
   IntStatus oldLevel = interrupt->SetLevel(IntOff);
   ASSERT(this == currentThread);
   DEBUG('t', "Yielding thread \"%s\"\n", getName());
   nextThread = scheduler->FindNextToRun();
   if (nextThread != NULL) {
scheduler->ReadyToRun(this);
scheduler->Run(nextThread);
    (void) interrupt->SetLevel(oldLevel);
// Thread::Sleep
// Relinquish the CPU, because the current thread is blocked
```

```
//waiting on a synchronization variable (Semaphore, Lock, or
Condition).
//Eventually, some thread will wake this thread up, and put it
//back on the ready queue, so that it can be re-scheduled.
//NOTE: if there are no threads on the ready queue, that means
//we have no thread to run. "Interrupt::Idle" is called
//to signify that we should idle the CPU until the next I/O interrupt
//occurs (the only thing that could cause a thread to become
//ready to run).
//NOTE: we assume interrupts are already disabled, because it
//is called from the synchronization routines which must
//disable interrupts for atomicity. We need interrupts off
//so that there can't be a time slice between pulling the first thread
//off the ready list, and switching to it.
//-----
biov
Thread::Sleep ()
   Thread *nextThread;
   ASSERT(this == currentThread);
   ASSERT(interrupt->getLevel() == IntOff);
   DEBUG('t', "Sleeping thread \"%s\"\n", getName());
   status = BLOCKED;
   while ((nextThread = scheduler->FindNextToRun()) == NULL)
interrupt->Idle();// no one to run, wait for an interrupt
   scheduler->Run(nextThread); // returns when we've been signalled
// ThreadFinish, InterruptEnable, ThreadPrint
//Dummy functions because C++ does not allow a pointer to a member
//function. So in order to do this, we create a dummy C function
//(which we can pass a pointer to), that then simply calls the
//member function.
static void ThreadFinish() { currentThread->Finish(); }
static void InterruptEnable() { interrupt->Enable(); }
void ThreadPrint(int arg){ Thread *t = (Thread *)arg; t->Print(); }
// Thread::StackAllocate
//Allocate and initialize an execution stack. The stack is
//initialized with an initial stack frame for ThreadRoot, which:
//enables interrupts
//calls (*func)(arg)
//calls Thread::Finish
//"func" is the procedure to be forked
//"arg" is the parameter to be passed to the procedure
Thread::StackAllocate (VoidFunctionPtr func, int arg)
```

```
stack = (int *) AllocBoundedArray(StackSize * sizeof(int));
#ifdef HOST SNAKE
   // HP stack works from low addresses to high addresses
   stackTop = stack + 16;// HP requires 64-byte frame marker
   stack[StackSize - 1] = STACK FENCEPOST;
#else
   // i386 & MIPS & SPARC stack works from high addresses to low
addresses
#ifdef HOST SPARC
   // SPARC stack must contains at least 1 activation record to start
with.
   stackTop = stack + StackSize - 96;
#else // HOST_MIPS || HOST_i386
   stackTop = stack + StackSize - 4i// - 4 to be on the safe side!
#ifdef HOST i386
   // the 80386 passes the return address on the stack. In order for
   // SWITCH() to go to ThreadRoot when we switch to this thread, the
   // return addres used in SWITCH() must be the starting address of
   // ThreadRoot.
   *(--stackTop) = (int)ThreadRoot;
#endif
#endif // HOST SPARC
  *stack = STACK FENCEPOST;
#endif // HOST_SNAKE
   machineState[PCState] = (int) ThreadRoot;
   machineState[StartupPCState] = (int) InterruptEnable;
   machineState[InitialPCState] = (int) func;
   machineState[InitialArgState] = arg;
   machineState[WhenDonePCState] = (int) ThreadFinish;
#ifdef USER PROGRAM
#include "machine.h"
//-----
// Thread::SaveUserState
//Save the CPU state of a user program on a context switch.
//Note that a user program thread has *two* sets of CPU registers --
//one for its state while executing user code, one for its state
//while executing kernel code. This routine saves the former.
//-----
void
Thread::SaveUserState()
   for (int i = 0; i < NumTotalRegs; i++)</pre>
userRegisters[i] = machine->ReadRegister(i);
//-----
// Thread::RestoreUserState
//Restore the CPU state of a user program on a context switch.
//Note that a user program thread has *two* sets of CPU registers --
//one for its state while executing user code, one for its state
//while executing kernel code. This routine restores the former.
//-----
void
```

```
Thread::RestoreUserState()
    for (int i = 0; i < NumTotalRegs; i++)</pre>
machine->WriteRegister(i, userRegisters[i]);
#endif
```

```
// thread.h
//Data structures for managing threads. A thread represents
//sequential execution of code within a program.
//So the state of a thread includes the program counter,
//the processor registers, and the execution stack.
// Note that because we allocate a fixed size stack for each
//thread, it is possible to overflow the stack -- for instance,
//by recursing to too deep a level. The most common reason
//for this occuring is allocating large data structures
//on the stack. For instance, this will cause problems:
//
//void foo() { int buf[1000]; ...}
//Instead, you should allocate all data structures dynamically:
//
//void foo() { int *buf = new int[1000]; ...}
//
// Bad things happen if you overflow the stack, and in the worst
//case, the problem may not be caught explicitly. Instead,
//the only symptom may be bizarre segmentation faults. (Of course,
//other problems can cause seq faults, so that isn't a sure sign
//that your thread stacks are too small.)
//One thing to try if you find yourself with seg faults is to
//increase the size of thread stack -- ThreadStackSize.
// In this interface, forking a thread takes two steps.
//We must first allocate a data structure for it: "t = new Thread".
//Only then can we do the fork: "t->fork(f, arg)".
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#ifndef THREAD_H
#define THREAD H
#include "copyright.h"
#include "utility.h"
#ifdef USER_PROGRAM
#include "machine.h"
#include "addrspace.h"
#endif
// CPU register state to be saved on context switch.
// The SPARC and MIPS only need 10 registers, but the Snake needs 18.
// For simplicity, this is just the max over all architectures.
#define MachineStateSize 18
// Size of the thread's private execution stack.
// WATCH OUT IF THIS ISN'T BIG ENOUGH!!!!!
#define StackSize(4 * 1024)// in words
// Thread state
enum ThreadStatus { JUST_CREATED, RUNNING, READY, BLOCKED };
// external function, dummy routine whose sole job is to call
Thread::Print
extern void ThreadPrint(int arg);
```

```
// The following class defines a "thread control block" -- which
// represents a single thread of execution.
//
   Every thread has:
       an execution stack for activation records ("stackTop" and
//
"stack")
       space to save CPU registers while not running ("machineState")
//
       a "status" (running/ready/blocked)
11
// Some threads also belong to a user address space; threads
   that only run in the kernel have a NULL address space.
class Thread {
 private:
    // NOTE: DO NOT CHANGE the order of these first two members.
    // THEY MUST be in this position for SWITCH to work.
    int* stackTop; // the current stack pointer
    int machineState[MachineStateSize]; // all registers except for
stackTop
 public:
    Thread(char* debugName);// initialize a Thread
    ~Thread(); // deallocate a Thread
// NOTE -- thread being deleted
// must not be running when delete
// is called
    // basic thread operations
    void Fork(VoidFunctionPtr func, int arg); // Make thread run
(*func)(arg)
    void Yield(); // Relinquish the CPU if any
// other thread is runnable
    void Sleep(); // Put the thread to sleep and
// relinguish the processor
    void Finish(); // The thread is done executing
    void CheckOverflow();
                            // Check if thread has
// overflowed its stack
    void setStatus(ThreadStatus st) { status = st; }
    char* getName() { return (name); }
    void Print() { printf("%s, ", name); }
 private:
    // some of the private data for this class is listed above
    int* stack; // Bottom of the stack
// NULL if this is the main thread
// (If NULL, don't deallocate stack)
    ThreadStatus status;// ready, running or blocked
    char* name;
    void StackAllocate(VoidFunctionPtr func, int arg);
    // Allocate a stack for thread.
// Used internally by Fork()
#ifdef USER PROGRAM
// A thread running a user program actually has *two* sets of CPU
registers --
// one for its state while executing user code, one for its state
// while executing kernel code.
    int userRegisters[NumTotalRegs];// user-level CPU register state
  public:
```

```
void SaveUserState();// save user-level register state
  void RestoreUserState();// restore user-level register state
  AddrSpace *space;// User code this thread is running.
#endif
};
// Magical machine-dependent routines, defined in switch.s
extern "C" {
// First frame on thread execution stack;
// enable interrupts
//call "func"
//(when func returns, if ever) call ThreadFinish()
void ThreadRoot();
// Stop running oldThread and start running newThread
void SWITCH(Thread *oldThread, Thread *newThread);
}
#endif // THREAD_H
```

```
// threadtest.cc
//Simple test case for the threads assignment.
//Create two threads, and have them context switch
//back and forth between themselves by calling Thread::Yield.
//to illustrate the inner workings of the thread system.
// Copyright (c) 1992-1993 The Regents of the University of California.
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limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "system.h"
// testnum is set in main.cc
int testnum = 1;
// SimpleThread
// Loop 5 times, yielding the CPU to another ready thread
//each iteration.
//"which" is simply a number identifying the thread, for debugging
//purposes.
void
SimpleThread(int which)
   int num;
   for (num = 0; num < 5; num++) {
printf("*** thread %d looped %d times\n", which, num);
       currentThread->Yield();
       _____
// ThreadTest1
// Set up a ping-pong between two threads, by forking a thread
//to call SimpleThread, and then calling SimpleThread ourselves.
void
ThreadTest1()
   DEBUG('t', "Entering ThreadTest1");
   Thread *t = new Thread("forked thread");
   t->Fork(SimpleThread, 1);
   SimpleThread(0);
               ._____
// ThreadTest
// Invoke a test routine.
void
ThreadTest()
   switch (testnum) {
   case 1:
ThreadTest1();
```

```
break;
    default:
printf("No test specified.\n");
break;
```

```
// utility.cc
//Debugging routines. Allows users to control whether to
//print DEBUG statements, based on a command line argument.
// Copyright (c) 1992-1993 The Regents of the University of California. va list ap;
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "utility.h"
// this seems to be dependent on how the compiler is configured.
// if you have problems with va start, try both of these alternatives
#ifdef HOST SNAKE
#include <stdarq.h>
#else
#ifdef HOST SPARC
#include <stdarq.h>
#else
#include "stdarg.h"
#endif
#endif
static char *enableFlags = NULL; // controls which DEBUG messages are
// DebugInit
      Initialize so that only DEBUG messages with a flag in flagList
//will be printed.
//If the flag is "+", we enable all DEBUG messages.
// "flagList" is a string of characters for whose DEBUG messages are
//to be enabled.
//----
void
DebugInit(char *flagList)
   enableFlags = flagList;
                  -----
// DebugIsEnabled
     Return TRUE if DEBUG messages with "flag" are to be printed.
//-----
bool
DebugIsEnabled(char flag)
   if (enableFlags != NULL)
     return (strchr(enableFlags, flag) != 0)
|| (strchr(enableFlags, '+') != 0);
   else
    return FALSE;
//-----
// DEBUG
      Print a debug message, if flag is enabled. Like printf,
//only with an extra argument on the front.
//-----
```

```
void
DEBUG(char flag, char *format, ...)
{
    if (DebugIsEnabled(flag)) {
    va_list ap;
    // You will get an unused variable message here -- ignore it.
    va_start(ap, format);
    vfprintf(stdout, format, ap);
    va_end(ap);
    fflush(stdout);
    }
}
```

```
// utility.h
//Miscellaneous useful definitions, including debugging routines.
//The debugging routines allow the user to turn on selected
//debugging messages, controllable from the command line arguments
//passed to Nachos (-d). You are encouraged to add your own
//debugging flags. The pre-defined debugging flags are:
//
//'+' -- turn on all debug messages
    't' -- thread system
    's' -- semaphores, locks, and conditions
     'i' -- interrupt emulation
//
     'm' -- machine emulation (USER PROGRAM)
//
     'd' -- disk emulation (FILESYS)
//
    'f' -- file system (FILESYS)
     'a' -- address spaces (USER PROGRAM)
//
     'n' -- network emulation (NETWORK)
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
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// of liability and disclaimer of warranty provisions.
#ifndef UTILITY H
#define UTILITY H
#include "copyright.h"
// Miscellaneous useful routines
#include "bool.h"
// Boolean values.
// This is the same definition
// as in the q++ library.
\#define min(a,b) (((a) < (b)) ? (a) : (b))
\#define \max(a,b) (((a) > (b)) ? (a) : (b))
// Divide and either round up or down
#define divRoundDown(n,s) ((n) / (s))
#define divRoundUp(n,s)
                          (((n) / (s)) + ((((n) % (s)) > 0) ? 1 : 0))
// This declares the type "VoidFunctionPtr" to be a "pointer to a
// function taking an integer argument and returning nothing". With
// such a function pointer (say it is "func"), we can call it like
this:
//
//(*func) (17);
// This is used by Thread::Fork and for interrupt handlers, as well
// as a couple of other places.
typedef void (*VoidFunctionPtr)(int arg);
typedef void (*VoidNoArgFunctionPtr)();
// Include interface that isolates us from the host machine system
library.
// Requires definition of bool, and VoidFunctionPtr
#include "sysdep.h"
// Interface to debugging routines.
extern void DebugInit(char* flags);// enable printing debug messages
extern bool DebugIsEnabled(char flag); // Is this debug flag enabled?
extern void DEBUG (char flag, char* format, ...); // Print debug
message
// if flag is enabled
```

```
// directory.cc
//Routines to manage a directory of file names.
//The directory is a table of fixed length entries; each
//entry represents a single file, and contains the file name.
//and the location of the file header on disk. The fixed size
//of each directory entry means that we have the restriction
//of a fixed maximum size for file names.
//The constructor initializes an empty directory of a certain size;
//we use ReadFrom/WriteBack to fetch the contents of the directory
//from disk, and to write back any modifications back to disk.
//
//Also, this implementation has the restriction that the size
//of the directory cannot expand. In other words, once all the
//entries in the directory are used, no more files can be created.
//Fixing this is one of the parts to the assignment.
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "utility.h"
#include "filehdr.h"
#include "directory.h"
//-----
// Directory::Directory
// Initialize a directory; initially, the directory is completely
//empty. If the disk is being formatted, an empty directory
//is all we need, but otherwise, we need to call FetchFrom in order
//to initialize it from disk.
//"size" is the number of entries in the directory
Directory::Directory(int size)
   table = new DirectoryEntry[size];
   tableSize = size;
   for (int i = 0; i < tableSize; i++)</pre>
table[i].inUse = FALSE;
´//-----
// Directory::~Directory
// De-allocate directory data structure.
//-----
Directory::~Directory()
   delete [] table;
.
//-----
// Directory::FetchFrom
// Read the contents of the directory from disk.
//
//"file" -- file containing the directory contents
//----
```

```
void
Directory::FetchFrom(OpenFile *file)
    (void) file->ReadAt((char *)table, tableSize *
sizeof(DirectoryEntry), 0);
// Directory::WriteBack
// Write any modifications to the directory back to disk
//"file" -- file to contain the new directory contents
biov
Directory::WriteBack(OpenFile *file)
    (void) file->WriteAt((char *)table, tableSize *
sizeof(DirectoryEntry), 0);
// Directory::FindIndex
// Look up file name in directory, and return its location in the
table of
//directory entries. Return -1 if the name isn't in the directory.
//
//"name" -- the file name to look up
Directory::FindIndex(char *name)
   for (int i = 0; i < tableSize; i++)</pre>
       if (table[i].inUse && !strncmp(table[i].name, name,
FileNameMaxLen))
   return i;
   return -1;// name not in directory
7/-----
// Directory::Find
// Look up file name in directory, and return the disk sector number
//where the file's header is stored. Return -1 if the name isn't
//in the directory.
//"name" -- the file name to look up
Directory::Find(char *name)
   int i = FindIndex(name);
   if (i != -1)
return table[i].sector;
   return -1;
7/----
// Directory::Add
// Add a file into the directory. Return TRUE if successful;
//return FALSE if the file name is already in the directory, or if
//the directory is completely full, and has no more space for
```

```
//additional file names.
//"name" -- the name of the file being added
//"newSector" -- the disk sector containing the added file's header
Directory::Add(char *name, int newSector)
   if (FindIndex(name) != -1)
return FALSE;
   for (int i = 0; i < tableSize; i++)
      if (!table[i].inUse) {
          table[i].inUse = TRUE;
          strncpy(table[i].name, name, FileNameMaxLen);
          table[i].sector = newSector;
      return TRUE;
   return FALSE; // no space. Fix when we have extensible files.
///-----
// Directory::Remove
// Remove a file name from the directory. Return TRUE if successful;
//return FALSE if the file isn't in the directory.
//
//"name" -- the file name to be removed
bool
Directory::Remove(char *name)
   int i = FindIndex(name);
   if (i == -1)
return FALSE; // name not in directory
   table[i].inUse = FALSE;
   return TRUE;
í/-----
// Directory::List
// List all the file names in the directory.
//-----
void
Directory::List()
  for (int i = 0; i < tableSize; i++)
if (table[i].inUse)
   printf("%s\n", table[i].name);
                _____
// Directory::Print
// List all the file names in the directory, their FileHeader
//and the contents of each file. For debugging.
void
Directory::Print()
   FileHeader *hdr = new FileHeader;
```

```
2/2
    printf("Directory contents:\n");
    for (int i = 0; i < tableSize; i++)</pre>
if (table[i].inUse) {
    printf("Name: %s, Sector: %d\n", table[i].name, table[i].sector);
   hdr->FetchFrom(table[i].sector);
   hdr->Print();
   printf("\n");
   delete hdr;
```

```
// directorv.h
//Data structures to manage a UNIX-like directory of file names.
//
       A directory is a table of pairs: <file name, sector #>,
//giving the name of each file in the directory, and
//where to find its file header (the data structure describing
//where to find the file's data blocks) on disk.
//
//
        We assume mutual exclusion is provided by the caller.
//
// Copyright (c) 1992-1993 The Regents of the University of California. // <file name, file header location>
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#ifndef DIRECTORY H
#define DIRECTORY H
#include "openfile.h"
#define FileNameMaxLen 9// for simplicity, we assume
// file names are <= 9 characters long
// The following class defines a "directory entry", representing a file
// in the directory. Each entry gives the name of the file, and where
// the file's header is to be found on disk.
//
// Internal data structures kept public so that Directory operations
// access them directly.
class DirectoryEntry {
 public:
   bool inUse;// Is this directory entry in use?
    int sector; // Location on disk to find the
// FileHeader for this file
   char name[FileNameMaxLen + 1];// Text name for file, with +1 for
// the trailing '\0'
};
// The following class defines a UNIX-like "directory". Each entry in
// the directory describes a file, and where to find it on disk.
//
// The directory data structure can be stored in memory, or on disk.
// When it is on disk, it is stored as a regular Nachos file.
//
// The constructor initializes a directory structure in memory; the
// FetchFrom/WriteBack operations shuffle the directory information
// from/to disk.
class Directory {
 public:
   Directory(int size); // Initialize an empty directory
// with space for "size" files
   ~Directory();// De-allocate the directory
   void FetchFrom(OpenFile *file); // Init directory contents from
disk
   void WriteBack(OpenFile *file);// Write modifications to
// directory contents back to disk
    int Find(char *name); // Find the sector number of the
// FileHeader for file: "name"
   bool Add(char *name, int newSector); // Add a file name into the
```

```
directory
   bool Remove(char *name); // Remove a file from the directory
   void List();// Print the names of all the files
// in the directory
   void Print();// Verbose print of the contents
// of the directory -- all the file
// names and their contents.
 private:
   int tableSize;// Number of directory entries
   DirectoryEntry *table;// Table of pairs:
   int FindIndex(char *name);// Find the index into the directory
// table corresponding to "name"
#endif // DIRECTORY H
```

```
// filehdr.cc
//Routines for managing the disk file header (in UNIX, this
//would be called the i-node).
//The file header is used to locate where on disk the
//file's data is stored. We implement this as a fixed size
//table of pointers -- each entry in the table points to the
//disk sector containing that portion of the file data
//(in other words, there are no indirect or doubly indirect
//blocks). The table size is chosen so that the file header
//will be just big enough to fit in one disk sector,
//
//
       Unlike in a real system, we do not keep track of file
permissions.
//ownership, last modification date, etc., in the file header.
//
//A file header can be initialized in two ways:
    for a new file, by modifying the in-memory data structure
//
      to point to the newly allocated data blocks
    for a file already on disk, by reading the file header from disk
//
// Copyright (c) 1992-1993 The Regents of the University of California. //
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "system.h"
#include "filehdr.h"
//-----
// FileHeader::Allocate
// Initialize a fresh file header for a newly created file.
//Allocate data blocks for the file out of the map of free disk blocks.
//Return FALSE if there are not enough free blocks to accomodate
//the new file.
//
//"freeMap" is the bit map of free disk sectors
//"fileSize" is the bit map of free disk sectors
//-----
FileHeader::Allocate(BitMap *freeMap, int fileSize)
   numBytes = fileSize;
   numSectors = divRoundUp(fileSize, SectorSize);
   if (freeMap->NumClear() < numSectors)</pre>
return FALSE; // not enough space
   for (int i = 0; i < numSectors; i++)</pre>
dataSectors[i] = freeMap->Find();
   return TRUE;
// FileHeader::Deallocate
// De-allocate all the space allocated for data blocks for this file.
//"freeMap" is the bit map of free disk sectors
void
```

```
FileHeader::Deallocate(BitMap *freeMap)
   for (int i = 0; i < numSectors; i++) {
ASSERT(freeMap->Test((int) dataSectors[i])); // ought to be marked!
freeMap->Clear((int) dataSectors[i]);
//-----
// FileHeader::FetchFrom
// Fetch contents of file header from disk.
//"sector" is the disk sector containing the file header
FileHeader::FetchFrom(int sector)
   synchDisk->ReadSector(sector, (char *)this);
// FileHeader::WriteBack
// Write the modified contents of the file header back to disk.
//"sector" is the disk sector to contain the file header
void
|FileHeader::WriteBack(int sector)
   synchDisk->WriteSector(sector, (char *)this);
//----
// FileHeader::ByteToSector
// Return which disk sector is storing a particular byte within the
file.
      This is essentially a translation from a virtual address (the
//offset in the file) to a physical address (the sector where the
//data at the offset is stored).
//"offset" is the location within the file of the byte in question
int
FileHeader::ByteToSector(int offset)
   return(dataSectors[offset / SectorSize]);
// FileHeader::FileLength
// Return the number of bytes in the file.
|FileHeader::FileLength()
   return numBytes;
//-----
// FileHeader::Print
// Print the contents of the file header, and the contents of all
```

```
//the data blocks pointed to by the file header.
void
FileHeader::Print()
    int i, j, k;
    char *data = new char[SectorSize];
    printf("FileHeader contents. File size: %d. File blocks:\n",
numBytes);
    for (i = 0; i < numSectors; i++)</pre>
printf("%d ", dataSectors[i]);
    printf("\nFile contents:\n");
    for (i = k = 0; i < numSectors; i++) {
synchDisk->ReadSector(dataSectors[i], data);
        for (j = 0; (j < SectorSize) && (k < numBytes); j++, k++) {}
    if ('\040' <= data[j] && data[j] <= '\176')  // isprint(data[j])</pre>
printf("%c", data[j]);
            else
printf("\\%x", (unsigned char)data[j]);
        printf("\n");
    delete [] data;
```

```
// filehdr.h
//Data structures for managing a disk file header.
//A file header describes where on disk to find the data in a file,
//along with other information about the file (for instance, its
//length, owner, etc.)
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#ifndef FILEHDR H
#define FILEHDR H
#include "disk.h"
#include "bitmap.h"
#define NumDirect ((SectorSize - 2 * sizeof(int)) / sizeof(int))
#define MaxFileSize (NumDirect * SectorSize)
// The following class defines the Nachos "file header" (in UNIX terms,
// the "i-node"), describing where on disk to find all of the data in
the file.
// The file header is organized as a simple table of pointers to
// data blocks.
//
// The file header data structure can be stored in memory or on disk.
// When it is on disk, it is stored in a single sector -- this means
// that we assume the size of this data structure to be the same
// as one disk sector. Without indirect addressing, this
// limits the maximum file length to just under 4K bytes.
//
// There is no constructor; rather the file header can be initialized
// by allocating blocks for the file (if it is a new file), or by
// reading it from disk.
class FileHeader {
 public:
   bool Allocate(BitMap *bitMap, int fileSize);// Initialize a file
header.
// including allocating space
// on disk for the file data
   void Deallocate(BitMap *bitMap); // De-allocate this file's
// data blocks
   void FetchFrom(int sectorNumber); // Initialize file header from
disk
   void WriteBack(int sectorNumber); // Write modifications to file
header
// back to disk
    int ByteToSector(int offset);// Convert a byte offset into the file
// to the disk sector containing
// the byte
    int FileLength();// Return the length of the file
// in bytes
   void Print();// Print the contents of the file.
 private:
    int numBytes; // Number of bytes in the file
    int numSectors; // Number of data sectors in the file
    int dataSectors[NumDirect]; // Disk sector numbers for each data
```

// block in the file #endif // FILEHDR H

```
// filesys.cc
//Routines to manage the overall operation of the file system.
//Implements routines to map from textual file names to files.
//
//Each file in the file system has:
     A file header, stored in a sector on disk
//(the size of the file header data structure is arranged
//to be precisely the size of 1 disk sector)
    A number of data blocks
    An entry in the file system directory
//
//
// The file system consists of several data structures:
    A bitmap of free disk sectors (cf. bitmap.h)
//
    A directory of file names and file headers
//
        Both the bitmap and the directory are represented as normal
//
//files. Their file headers are located in specific sectors
//(sector 0 and sector 1), so that the file system can find them
//on bootup.
//
//The file system assumes that the bitmap and directory files are
//kept "open" continuously while Nachos is running.
//For those operations (such as Create, Remove) that modify the
//directory and/or bitmap, if the operation succeeds, the changes
//are written immediately back to disk (the two files are kept
//open during all this time). If the operation fails, and we have
//modified part of the directory and/or bitmap, we simply discard
//the changed version, without writing it back to disk.
//
// Our implementation at this point has the following restrictions:
//
//
     there is no synchronization for concurrent accesses
     files have a fixed size, set when the file is created
//
     files cannot be bigger than about 3KB in size
//
//
     there is no hierarchical directory structure, and only a limited
//
       number of files can be added to the system
     there is no attempt to make the system robust to failures
//
//
     (if Nachos exits in the middle of an operation that modifies
      the file system, it may corrupt the disk)
//
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "disk.h"
#include "bitmap.h"
#include "directory.h"
#include "filehdr.h"
#include "filesys.h"
// Sectors containing the file headers for the bitmap of free sectors,
// and the directory of files. These file headers are placed in well-
known
// sectors, so that they can be located on boot-up.
#define FreeMapSector 0
```

```
#define DirectorySector 1
// Initial file sizes for the bitmap and directory; until the file
// supports extensible files, the directory size sets the maximum
number
// of files that can be loaded onto the disk.
#define FreeMapFileSize (NumSectors / BitsInByte)
#define NumDirEntries 10
#define DirectoryFileSize (sizeof(DirectoryEntry) * NumDirEntries)
// FileSystem::FileSystem
// Initialize the file system. If format = TRUE, the disk has
//nothing on it, and we need to initialize the disk to contain
//an empty directory, and a bitmap of free sectors (with almost but
//not all of the sectors marked as free).
//
//If format = FALSE, we just have to open the files
//representing the bitmap and the directory.
//
//"format" -- should we initialize the disk?
FileSystem::FileSystem(bool format)
    DEBUG('f', "Initializing the file system.\n");
    if (format) {
        BitMap *freeMap = new BitMap(NumSectors);
        Directory *directory = new Directory(NumDirEntries);
FileHeader *mapHdr = new FileHeader;
FileHeader *dirHdr = new FileHeader;
        DEBUG('f', "Formatting the file system.\n");
    // First, allocate space for FileHeaders for the directory and
bitmap
    // (make sure no one else grabs these!)
freeMap->Mark(FreeMapSector);
freeMap->Mark(DirectorySector);
    // Second, allocate space for the data blocks containing the
contents
    // of the directory and bitmap files. There better be enough
space!
ASSERT(mapHdr->Allocate(freeMap, FreeMapFileSize));
ASSERT(dirHdr->Allocate(freeMap, DirectoryFileSize));
    // Flush the bitmap and directory FileHeaders back to disk
    // We need to do this before we can "Open" the file, since open
    // reads the file header off of disk (and currently the disk has
garbage
   // on it!).
        DEBUG('f', "Writing headers back to disk.\n");
mapHdr->WriteBack(FreeMapSector);
dirHdr->WriteBack(DirectorySector);
    // OK to open the bitmap and directory files now
    // The file system operations assume these two files are left open
    // while Nachos is running.
        freeMapFile = new OpenFile(FreeMapSector);
        directoryFile = new OpenFile(DirectorySector);
    // Once we have the files "open", we can write the initial version
```

```
// of each file back to disk. The directory at this point is
completely
   // empty; but the bitmap has been changed to reflect the fact that
   // sectors on the disk have been allocated for the file headers and
   // to hold the file data for the directory and bitmap.
       DEBUG('f', "Writing bitmap and directory back to disk.\n");
freeMap->WriteBack(freeMapFile); // flush changes to disk
directory->WriteBack(directoryFile);
if (DebugIsEnabled('f')) {
   freeMap->Print();
   directory->Print();
       delete freeMap;
delete directory;
delete mapHdr;
delete dirHdr;
   } else {
   // if we are not formatting the disk, just open the files
representing
   // the bitmap and directory; these are left open while Nachos is
running
       freeMapFile = new OpenFile(FreeMapSector);
       directoryFile = new OpenFile(DirectorySector);
//----
// FileSystem::Create
// Create a file in the Nachos file system (similar to UNIX create).
//Since we can't increase the size of files dynamically, we have
//to give Create the initial size of the file.
//
//The steps to create a file are:
// Make sure the file doesn't already exist
         Allocate a sector for the file header
// Allocate space on disk for the data blocks for the file
// Add the name to the directory
// Store the new file header on disk
// Flush the changes to the bitmap and the directory back to disk
//Return TRUE if everything goes ok, otherwise, return FALSE.
//
// Create fails if:
// file is already in directory
// no free space for file header
// no free entry for file in directory
// no free space for data blocks for the file
// Note that this implementation assumes there is no concurrent access |FileSystem::Open(char *name)
//to the file system!
//"name" -- name of file to be created
//"initialSize" -- size of file to be created
FileSystem::Create(char *name, int initialSize)
```

```
Directory *directory;
   BitMap *freeMap;
   FileHeader *hdr;
   int sector;
   bool success;
   DEBUG('f', "Creating file %s, size %d\n", name, initialSize);
   directory = new Directory(NumDirEntries);
   directory->FetchFrom(directoryFile);
   if (directory->Find(name) != -1)
      success = FALSE;// file is already in directory
   else {
        freeMap = new BitMap(NumSectors);
        freeMap->FetchFrom(freeMapFile);
        sector = freeMap->Find();// find a sector to hold the file
header
   if (sector == -1)
           success = FALSE; // no free block for file header
       else if (!directory->Add(name, sector))
           success = FALSE;// no space in directory
else {
       hdr = new FileHeader;
   if (!hdr->Allocate(freeMap, initialSize))
           success = FALSE; // no space on disk for data
   else {
    success = TRUE;
// everthing worked, flush all changes back to disk
       hdr->WriteBack(sector);
       directory->WriteBack(directoryFile);
       freeMap->WriteBack(freeMapFile);
           delete hdr;
       delete freeMap;
   delete directory;
   return success;
//-----
// FileSystem::Open
// Open a file for reading and writing.
//To open a file:
// Find the location of the file's header, using the directory
// Bring the header into memory
//"name" -- the text name of the file to be opened
OpenFile *
   Directory *directory = new Directory(NumDirEntries);
   OpenFile *openFile = NULL;
   int sector;
   DEBUG('f', "Opening file %s\n", name);
   directory->FetchFrom(directoryFile);
    sector = directory->Find(name);
   if (sector >= 0)
```

```
openFile = new OpenFile(sector);// name was found in directory
   delete directory;
   return openFile; // return NULL if not found
,
//-----
// FileSystem::Remove
// Delete a file from the file system. This requires:
    Remove it from the directory
//
    Delete the space for its header
    Delete the space for its data blocks
//
    Write changes to directory, bitmap back to disk
//
//
//Return TRUE if the file was deleted, FALSE if the file wasn't
//in the file system.
//"name" -- the text name of the file to be removed
//-----
bool
FileSystem::Remove(char *name)
   Directory *directory;
   BitMap *freeMap;
   FileHeader *fileHdr;
   int sector;
   directory = new Directory(NumDirEntries);
   directory->FetchFrom(directoryFile);
   sector = directory->Find(name);
   if (sector == -1) {
     delete directory;
     return FALSE; // file not found
   fileHdr = new FileHeader;
   fileHdr->FetchFrom(sector);
   freeMap = new BitMap(NumSectors);
   freeMap->FetchFrom(freeMapFile);
   fileHdr->Deallocate(freeMap); // remove data blocks
   freeMap->Clear(sector);// remove header block
   directory->Remove(name);
   freeMap->WriteBack(freeMapFile);// flush to disk
   delete fileHdr;
   delete directory;
   delete freeMap;
   return TRUE;
// FileSystem::List
// List all the files in the file system directory.
//-----
void
FileSystem::List()
   Directory *directory = new Directory(NumDirEntries);
   directory->FetchFrom(directoryFile);
   directory->List();
```

```
delete directory;
ĺ/----
// FileSystem::Print
// Print everything about the file system:
// the contents of the bitmap
// the contents of the directory
// for each file in the directory.
      the contents of the file header
//
//
      the data in the file
//-----
FileSystem::Print()
   FileHeader *bitHdr = new FileHeader;
   FileHeader *dirHdr = new FileHeader;
   BitMap *freeMap = new BitMap(NumSectors);
   Directory *directory = new Directory(NumDirEntries);
   printf("Bit map file header:\n");
   bitHdr->FetchFrom(FreeMapSector);
   bitHdr->Print();
   printf("Directory file header:\n");
   dirHdr->FetchFrom(DirectorySector);
   dirHdr->Print();
   freeMap->FetchFrom(freeMapFile);
   freeMap->Print();
   directory->FetchFrom(directoryFile);
   directory->Print();
   delete bitHdr;
   delete dirHdr;
   delete freeMap;
   delete directory;
```

```
// filesys.h
//Data structures to represent the Nachos file system.
//A file system is a set of files stored on disk, organized
//into directories. Operations on the file system have to
//do with "naming" -- creating, opening, and deleting files,
//qiven a textual file name. Operations on an individual
//"open" file (read, write, close) are to be found in the OpenFile
//class (openfile.h).
//
//We define two separate implementations of the file system.
//The "STUB" version just re-defines the Nachos file system
//operations as operations on the native UNIX file system on the
machine
//running the Nachos simulation. This is provided in case the
//multiprogramming and virtual memory assignments (which make use
//of the file system) are done before the file system assignment.
//The other version is a "real" file system, built on top of
//a disk simulator. The disk is simulated using the native UNIX
//file system (in a file named "DISK").
//In the "real" implementation, there are two key data structures used
//in the file system. There is a single "root" directory, listing
//all of the files in the file system; unlike UNIX, the baseline
//system does not provide a hierarchical directory structure.
//In addition, there is a bitmap for allocating
//disk sectors. Both the root directory and the bitmap are themselves
//stored as files in the Nachos file system -- this causes an
interesting
//bootstrap problem when the simulated disk is initialized.
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#ifndef FS_H
#define FS_H
#include "copyright.h"
#include "openfile.h"
#ifdef FILESYS STUB // Temporarily implement file system calls as
// calls to UNIX, until the real file system
// implementation is available
class FileSystem {
 public:
   FileSystem(bool format) {}
   bool Create(char *name, int initialSize) {
int fileDescriptor = OpenForWrite(name);
if (fileDescriptor == -1) return FALSE;
Close(fileDescriptor);
return TRUE;
    OpenFile* Open(char *name) {
  int fileDescriptor = OpenForReadWrite(name, FALSE);
 if (fileDescriptor == -1) return NULL;
 return new OpenFile(fileDescriptor);
```

```
bool Remove(char *name) { return Unlink(name) == 0; }
#else // FILESYS
class FileSystem {
  public:
    FileSystem(bool format);// Initialize the file system.
 // Must be called *after* "synchDisk"
 // has been initialized.
    // If "format", there is nothing on
 // the disk, so initialize the directory
    // and the bitmap of free blocks.
    bool Create(char *name, int initialSize);
 // Create a file (UNIX creat)
    OpenFile* Open(char *name); // Open a file (UNIX open)
    bool Remove(char *name); // Delete a file (UNIX unlink)
    void List();// List all the files in the file system
    void Print();// List all the files and their contents
  private:
   OpenFile* freeMapFile;// Bit map of free disk blocks,
 // represented as a file
   OpenFile* directoryFile;// "Root" directory -- list of
// file names, represented as a file
#endif // FILESYS
#endif // FS_H
```

```
// fstest.cc
//Simple test routines for the file system.
//We implement:
// Copy -- copy a file from UNIX to Nachos
// Print -- cat the contents of a Nachos file
// Perftest -- a stress test for the Nachos file system
//read and write a really large file in tiny chunks
//(won't work on baseline system!)
//
// Copyright (c) 1992-1993 The Regents of the University of California. // Print the contents of the Nachos file "name".
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "utility.h"
#include "filesys.h"
#include "system.h"
#include "thread.h"
#include "disk.h"
#include "stats.h"
#define TransferSize 10 // make it small, just to be difficult
//-----
// Copy
// Copy the contents of the UNIX file "from" to the Nachos file "to"
void
Copy(char *from, char *to)
   FILE *fp;
   OpenFile* openFile;
   int amountRead, fileLength;
   char *buffer;
// Open UNIX file
   if ((fp = fopen(from, "r")) == NULL) {
printf("Copy: couldn't open input file %s\n", from);
return;
// Figure out length of UNIX file
   fseek(fp, 0, 2);
   fileLength = ftell(fp);
   fseek(fp, 0, 0);
// Create a Nachos file of the same length
   DEBUG('f', "Copying file %s, size %d, to file %s\n", from,
fileLength, to);
   if (!fileSystem->Create(to, fileLength)) { // Create Nachos file
printf("Copy: couldn't create output file %s\n", to);
fclose(fp);
return;
   openFile = fileSystem->Open(to);
   ASSERT(openFile != NULL);
// Copy the data in TransferSize chunks
   buffer = new char[TransferSize];
```

```
while ((amountRead = fread(buffer, sizeof(char), TransferSize,
fp)) > 0
openFile->Write(buffer, amountRead);
   delete [] buffer;
// Close the UNIX and the Nachos files
   delete openFile;
   fclose(fp);
,
//-----
// Print
//-----
biov
Print(char *name)
   OpenFile *openFile;
   int i, amountRead;
   char *buffer;
   if ((openFile = fileSystem->Open(name)) == NULL) {
printf("Print: unable to open file %s\n", name);
return;
   buffer = new char[TransferSize];
   while ((amountRead = openFile->Read(buffer, TransferSize)) > 0)
for (i = 0; i < amountRead; i++)</pre>
   printf("%c", buffer[i]);
   delete [] buffer;
   delete openFile; // close the Nachos file
   return;
,
,
,
,
// PerformanceTest
// Stress the Nachos file system by creating a large file, writing
//it out a bit at a time, reading it back a bit at a time, and then
//deleting the file.
//
//Implemented as three separate routines:
// FileWrite -- write the file
// FileRead -- read the file
// PerformanceTest -- overall control, and print out performance #'s
//-----
#define FileName "TestFile"
#define Contents "1234567890"
#define ContentSize strlen(Contents)
#define FileSize ((int)(ContentSize * 5000))
static void
FileWrite()
   OpenFile *openFile;
   int i, numBytes;
   printf("Sequential write of %d byte file, in %d byte chunks\n",
FileSize, ContentSize);
   if (!fileSystem->Create(FileName, 0)) {
     printf("Perf test: can't create %s\n", FileName);
```

```
openFile = fileSystem->Open(FileName);
    if (openFile == NULL) {
printf("Perf test: unable to open %s\n", FileName);
return;
    for (i = 0; i < FileSize; i += ContentSize) {</pre>
        numBytes = openFile->Write(Contents, ContentSize);
if (numBytes < 10) {
    printf("Perf test: unable to write %s\n", FileName);
    delete openFile;
    return;
    delete openFile;// close file
static void
FileRead()
    OpenFile *openFile;
    char *buffer = new char[ContentSize];
    int i, numBytes;
    printf("Sequential read of %d byte file, in %d byte chunks\n",
FileSize, ContentSize);
    if ((openFile = fileSystem->Open(FileName)) == NULL) {
printf("Perf test: unable to open file %s\n", FileName);
delete [] buffer;
return;
    for (i = 0; i < FileSize; i += ContentSize) {
        numBytes = openFile->Read(buffer, ContentSize);
if ((numBytes < 10) || strncmp(buffer, Contents, ContentSize)) {</pre>
    printf("Perf test: unable to read %s\n", FileName);
    delete openFile;
    delete [] buffer;
    return;
    delete [] buffer;
    delete openFile;// close file
void
PerformanceTest()
    printf("Starting file system performance test:\n");
    stats->Print();
    FileWrite();
    FileRead();
    if (!fileSystem->Remove(FileName)) {
      printf("Perf test: unable to remove %s\n", FileName);
      return;
    stats->Print();
```

```
// openfile.cc
//Routines to manage an open Nachos file. As in UNIX, a
//file must be open before we can read or write to it.
//Once we're all done, we can close it (in Nachos, by deleting
//the OpenFile data structure).
//Also as in UNIX, for convenience, we keep the file header in
//memory while the file is open.
// Copyright (c) 1992-1993 The Regents of the University of California. int
// All rights reserved. See copyright.h for copyright notice and
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "filehdr.h"
#include "openfile.h"
#include "system.h"
#ifdef HOST_SPARC
#include <strings.h>
//-----
// OpenFile::OpenFile
// Open a Nachos file for reading and writing. Bring the file header
//into memory while the file is open.
//"sector" -- the location on disk of the file header for this file
//----
OpenFile::OpenFile(int sector)
   hdr = new FileHeader;
   hdr->FetchFrom(sector);
   seekPosition = 0;
´//-----
// OpenFile::~OpenFile
// Close a Nachos file, de-allocating any in-memory data structures.
OpenFile::~OpenFile()
   delete hdr;
.
// OpenFile::Seek
// Change the current location within the open file -- the point at
//which the next Read or Write will start from.
//
//"position" -- the location within the file for the next Read/Write
void
OpenFile::Seek(int position)
   seekPosition = position;
//----
// OpenFile::Read/Write
// Read/write a portion of a file, starting from seekPosition.
```

```
//Return the number of bytes actually written or read, and as a
//side effect, increment the current position within the file.
//Implemented using the more primitive ReadAt/WriteAt.
//"into" -- the buffer to contain the data to be read from disk
//"from" -- the buffer containing the data to be written to disk
//"numBytes" -- the number of bytes to transfer
//-----
OpenFile::Read(char *into, int numBytes)
  int result = ReadAt(into, numBytes, seekPosition);
  seekPosition += result;
  return result;
OpenFile::Write(char *into, int numBytes)
  int result = WriteAt(into, numBytes, seekPosition);
  seekPosition += result;
  return result;
//-----
// OpenFile::ReadAt/WriteAt
// Read/write a portion of a file, starting at "position".
//Return the number of bytes actually written or read, but has
//no side effects (except that Write modifies the file, of course).
//There is no guarantee the request starts or ends on an even disk
//boundary; however the disk only knows how to read/write a whole disk
//sector at a time. Thus:
//For ReadAt:
//
    We read in all of the full or partial sectors that are part of
the
    request, but we only copy the part we are interested in.
//
//For WriteAt:
    We must first read in any sectors that will be partially written,
    so that we don't overwrite the unmodified portion. We then copy
//
    in the data that will be modified, and write back all the full
    or partial sectors that are part of the request.
//
//
//"into" -- the buffer to contain the data to be read from disk
//"from" -- the buffer containing the data to be written to disk
//"numBytes" -- the number of bytes to transfer
//"position" -- the offset within the file of the first byte to be
//read/written
OpenFile::ReadAt(char *into, int numBytes, int position)
   int fileLength = hdr->FileLength();
   int i, firstSector, lastSector, numSectors;
   char *buf;
```

```
if ((numBytes <= 0) || (position >= fileLength))
    return 0; // check request
    if ((position + numBytes) > fileLength)
numBytes = fileLength - position;
    DEBUG('f', "Reading %d bytes at %d, from file of length %d.\n",
numBytes, position, fileLength);
    firstSector = divRoundDown(position, SectorSize);
    lastSector = divRoundDown(position + numBytes - 1, SectorSize);
    numSectors = 1 + lastSector - firstSector;
    // read in all the full and partial sectors that we need
    buf = new char[numSectors * SectorSize];
    for (i = firstSector; i <= lastSector; i++)</pre>
        synchDisk->ReadSector(hdr->ByteToSector(i * SectorSize),
&buf[(i - firstSector) * SectorSize]);
    // copy the part we want
    bcopy(&buf[position - (firstSector * SectorSize)], into, numBytes);
    delete [] buf;
    return numBytes;
int
OpenFile::WriteAt(char *from, int numBytes, int position)
    int fileLength = hdr->FileLength();
    int i, firstSector, lastSector, numSectors;
    bool firstAligned, lastAligned;
    char *buf;
    if ((numBytes <= 0) | (position >= fileLength))
return 0;// check request
    if ((position + numBytes) > fileLength)
numBytes = fileLength - position;
    DEBUG('f', "Writing %d bytes at %d, from file of length %d.\n",
numBytes, position, fileLength);
    firstSector = divRoundDown(position, SectorSize);
    lastSector = divRoundDown(position + numBytes - 1, SectorSize);
    numSectors = 1 + lastSector - firstSector;
    buf = new char[numSectors * SectorSize];
    firstAligned = (position == (firstSector * SectorSize));
    lastAligned = ((position + numBytes) == ((lastSector + 1) *
SectorSize));
// read in first and last sector, if they are to be partially modified
    if (!firstAligned)
        ReadAt(buf, SectorSize, firstSector * SectorSize);
    if (!lastAligned && ((firstSector != lastSector) | firstAligned))
        ReadAt(&buf[(lastSector - firstSector) * SectorSize],
SectorSize, lastSector * SectorSize);
// copy in the bytes we want to change
    bcopy(from, &buf[position - (firstSector * SectorSize)], numBytes);
// write modified sectors back
    for (i = firstSector; i <= lastSector; i++)</pre>
        synchDisk->WriteSector(hdr->ByteToSector(i * SectorSize),
&buf[(i - firstSector) * SectorSize]);
    delete [] buf;
    return numBytes;
// OpenFile::Length
```

```
// Return the number of bytes in the file.
OpenFile::Length()
   return hdr->FileLength();
```

```
// openfile.h
//Data structures for opening, closing, reading and writing to
//individual files. The operations supported are similar to
//the UNIX ones -- type 'man open' to the UNIX prompt.
//
//There are two implementations. One is a "STUB" that directly
//turns the file operations into the underlying UNIX operations.
//(cf. comment in filesys.h).
//The other is the "real" implementation, that turns these
//operations into read and write disk sector requests.
//In this baseline implementation of the file system, we don't
//worry about concurrent accesses to the file system
//by different threads -- this is part of the assignment.
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#ifndef OPENFILE H
#define OPENFILE H
#include "copyright.h"
#include "utility.h"
#ifdef FILESYS_STUB// Temporarily implement calls to
// Nachos file system as calls to UNIX!
// See definitions listed under #else
class OpenFile {
 public:
   OpenFile(int f) { file = f; currentOffset = 0; }// open the file
    ~OpenFile() { Close(file); }// close the file
    int ReadAt(char *into, int numBytes, int position) {
   Lseek(file, position, 0);
return ReadPartial(file, into, numBytes);
   int WriteAt(char *from, int numBytes, int position) {
   Lseek(file, position, 0);
WriteFile(file, from, numBytes);
return numBytes;
    int Read(char *into, int numBytes) {
int numRead = ReadAt(into, numBytes, currentOffset);
currentOffset += numRead;
return numRead;
    int Write(char *from, int numBytes) {
int numWritten = WriteAt(from, numBytes, currentOffset);
currentOffset += numWritten;
return numWritten;
   int Length() { Lseek(file, 0, 2); return Tell(file); }
 private:
   int file;
   int currentOffset;
#else // FILESYS
```

```
class FileHeader;
class OpenFile {
 public:
    OpenFile(int sector); // Open a file whose header is located
// at "sector" on the disk
    ~OpenFile();// Close the file
    void Seek(int position); // Set the position from which to
// start reading/writing -- UNIX lseek
    int Read(char *into, int numBytes); // Read/write bytes from the
file.
// starting at the implicit position.
// Return the # actually read/written,
// and increment position in file.
    int Write(char *from, int numBytes);
    int ReadAt(char *into, int numBytes, int position);
    // Read/write bytes from the file,
// bypassing the implicit position.
    int WriteAt(char *from, int numBytes, int position);
    int Length(); // Return the number of bytes in the
// file (this interface is simpler
// than the UNIX idiom -- lseek to
// end of file, tell, lseek back
 private:
    FileHeader *hdr; // Header for this file
    int seekPosition;// Current position within the file
#endif // FILESYS
#endif // OPENFILE H
```

```
// synchdisk.cc
//Routines to synchronously access the disk. The physical disk
//is an asynchronous device (disk requests return immediately, and
//an interrupt happens later on). This is a layer on top of
//the disk providing a synchronous interface (requests wait until
//the request completes).
//Use a semaphore to synchronize the interrupt handlers with the
//pending requests. And, because the physical disk can only
//handle one operation at a time, use a lock to enforce mutual
//exclusion.
#include "copyright.h"
#include "synchdisk.h"
//-----
// DiskRequestDone
// Disk interrupt handler. Need this to be a C routine, because
//C++ can't handle pointers to member functions.
//----
static void
DiskRequestDone (int arg)
   SynchDisk* disk = (SynchDisk *)arg;
   disk->RequestDone();
í/-----
// SynchDisk::SynchDisk
// Initialize the synchronous interface to the physical disk, in turn
//initializing the physical disk.
//"name" -- UNIX file name to be used as storage for the disk data
// (usually, "DISK")
//----
SynchDisk::SynchDisk(char* name)
   semaphore = new Semaphore("synch disk", 0);
   lock = new Lock("synch disk lock");
   disk = new Disk(name, DiskRequestDone, (int) this);
// SynchDisk::~SynchDisk
// De-allocate data structures needed for the synchronous disk
//abstraction.
//-----
SynchDisk::~SynchDisk()
   delete disk;
   delete lock;
   delete semaphore;
                _____
// SynchDisk::ReadSector
// Read the contents of a disk sector into a buffer. Return only
//after the data has been read.
//
//"sectorNumber" -- the disk sector to read
//"data" -- the buffer to hold the contents of the disk sector
```

```
biov
SynchDisk::ReadSector(int sectorNumber, char* data)
   lock->Acquire();// only one disk I/O at a time
   disk->ReadRequest(sectorNumber, data);
   semaphore->P();// wait for interrupt
   lock->Release();
    _____
// SynchDisk::WriteSector
// Write the contents of a buffer into a disk sector. Return only
//after the data has been written
//"sectorNumber" -- the disk sector to be written
//"data" -- the new contents of the disk sector
//-----
biov
SynchDisk::WriteSector(int sectorNumber, char* data)
   lock->Acquire();// only one disk I/O at a time
   disk->WriteRequest(sectorNumber, data);
   semaphore->P();// wait for interrupt
   lock->Release();
.
.
.
// SynchDisk::RequestDone
// Disk interrupt handler. Wake up any thread waiting for the disk
//request to finish.
//----
SynchDisk::RequestDone()
   semaphore->V();
```

```
// synchdisk.h
// Data structures to export a synchronous interface to the raw
//disk device.
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#ifndef SYNCHDISK H
#define SYNCHDISK H
#include "disk.h"
#include "synch.h"
// The following class defines a "synchronous" disk abstraction.
// As with other I/O devices, the raw physical disk is an asynchronous
device --
// requests to read or write portions of the disk return immediately,
// and an interrupt occurs later to signal that the operation
completed.
// (Also, the physical characteristics of the disk device assume that
// only one operation can be requested at a time).
// This class provides the abstraction that for any individual thread
// making a request, it waits around until the operation finishes
before
// returning.
class SynchDisk {
 public:
   SynchDisk(char* name);
                              // Initialize a synchronous disk,
// by initializing the raw Disk.
    ~SynchDisk();// De-allocate the synch disk data
   void ReadSector(int sectorNumber, char* data);
   // Read/write a disk sector, returning
   // only once the data is actually read
// or written. These call
   // Disk::ReadRequest/WriteRequest and
// then wait until the request is done.
   void WriteSector(int sectorNumber, char* data);
   void RequestDone();// Called by the disk device interrupt
// handler, to signal that the
// current disk operation is complete.
 private:
   Disk *disk; // Raw disk device
   Semaphore *semaphore; // To synchronize requesting thread
// with the interrupt handler
   Lock *lock; // Only one read/write request
// can be sent to the disk at a time
};
#endif // SYNCHDISK H
```

```
// addrspace.cc
//Routines to manage address spaces (executing user programs).
//In order to run a user program, you must:
//
//1. link with the -N -T 0 option
//2. run coff2noff to convert the object file to Nachos format
//(Nachos object code format is essentially just a simpler
//version of the UNIX executable object code format)
//3. load the NOFF file into the Nachos file system
//(if you haven't implemented the file system yet, you
//don't need to do this last step)
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "system.h"
#include "addrspace.h"
#include "noff.h"
#ifdef HOST SPARC
#include <strings.h>
#endif
//-----
// SwapHeader
// Do little endian to big endian conversion on the bytes in the
//object file header, in case the file was generated on a little
//endian machine, and we're now running on a big endian machine.
//-----
static void
SwapHeader (NoffHeader *noffH)
noffH->noffMagic = WordToHost(noffH->noffMagic);
noffH->code.size = WordToHost(noffH->code.size);
noffH->code.virtualAddr = WordToHost(noffH->code.virtualAddr);
noffH->code.inFileAddr = WordToHost(noffH->code.inFileAddr);
noffH->initData.size = WordToHost(noffH->initData.size);
noffH->initData.virtualAddr = WordToHost(noffH->initData.virtualAddr);
noffH->initData.inFileAddr = WordToHost(noffH->initData.inFileAddr);
noffH->uninitData.size = WordToHost(noffH->uninitData.size);
noffH->uninitData.virtualAddr = WordToHost(noffH-
>uninitData.virtualAddr);
noffH->uninitData.inFileAddr = WordToHost(noffH-
>uninitData.inFileAddr);
// AddrSpace::AddrSpace
// Create an address space to run a user program.
//Load the program from a file "executable", and set everything
//up so that we can start executing user instructions.
//Assumes that the object code file is in NOFF format.
//
//First, set up the translation from program memory to physical
//memory. For now, this is really simple (1:1), since we are
```

```
//only uniprogramming, and we have a single unsegmented page table
//"executable" is the file containing the object code to load into
AddrSpace::AddrSpace(OpenFile *executable)
    NoffHeader noffH;
   unsigned int i, size;
    executable->ReadAt((char *)&noffH, sizeof(noffH), 0);
    if ((noffH.noffMagic != NOFFMAGIC) &&
(WordToHost(noffH.noffMagic) == NOFFMAGIC))
    SwapHeader(&noffH);
    ASSERT(noffH.noffMagic == NOFFMAGIC);
// how big is address space?
    size = noffH.code.size + noffH.initData.size +
noffH.uninitData.size
+ UserStackSize; // we need to increase the size
// to leave room for the stack
    numPages = divRoundUp(size, PageSize);
    size = numPages * PageSize;
   ASSERT(numPages <= NumPhysPages);// check we're not trying
// to run anything too big --
// at least until we have
// virtual memory
    DEBUG('a', "Initializing address space, num pages %d, size %d\n",
numPages, size);
// first, set up the translation
    pageTable = new TranslationEntry[numPages];
    for (i = 0; i < numPages; i++) {
pageTable[i].virtualPage = i;// for now, virtual page # = phys page #
pageTable[i].physicalPage = i;
pageTable[i].valid = TRUE;
pageTable[i].use = FALSE;
pageTable[i].dirty = FALSE;
pageTable[i].readOnly = FALSE; // if the code segment was entirely on
// a separate page, we could set its
// pages to be read-only
// zero out the entire address space, to zero the unitialized data
segment
// and the stack segment
   bzero(machine->mainMemory, size);
// then, copy in the code and data segments into memory
   if (noffH.code.size > 0)
        DEBUG('a', "Initializing code segment, at 0x%x, size %d\n",
noffH.code.virtualAddr, noffH.code.size);
        executable->ReadAt(&(machine-
>mainMemory[noffH.code.virtualAddr]),
noffH.code.size, noffH.code.inFileAddr);
    if (noffH.initData.size > 0) {
        DEBUG('a', "Initializing data segment, at 0x%x, size %d\n",
noffH.initData.virtualAddr, noffH.initData.size);
        executable->ReadAt(&(machine-
```

```
>mainMemory[noffH.initData.virtualAddr]),
noffH.initData.size, noffH.initData.inFileAddr);
1/-----
// AddrSpace::~AddrSpace
// Dealloate an address space. Nothing for now!
//-----
AddrSpace::~AddrSpace()
  delete pageTable;
   ______
// AddrSpace::InitRegisters
// Set the initial values for the user-level register set.
//
// We write these directly into the "machine" registers, so
//that we can immediately jump to user code. Note that these
//will be saved/restored into the currentThread->userRegisters
//when this thread is context switched out.
void
AddrSpace::InitRegisters()
   int i;
   for (i = 0; i < NumTotalRegs; i++)
machine->WriteRegister(i, 0);
   // Initial program counter -- must be location of "Start"
   machine->WriteRegister(PCReg, 0);
   // Need to also tell MIPS where next instruction is, because
   // of branch delay possibility
   machine->WriteRegister(NextPCReg, 4);
  // Set the stack register to the end of the address space, where we
  // allocated the stack; but subtract off a bit, to make sure we
don't
  // accidentally reference off the end!
   machine->WriteRegister(StackReg, numPages * PageSize - 16);
   DEBUG('a', "Initializing stack register to %d\n", numPages *
PageSize - 16);
// AddrSpace::SaveState
// On a context switch, save any machine state, specific
//to this address space, that needs saving.
//For now, nothing!
void AddrSpace::SaveState()
    _____
// AddrSpace::RestoreState
// On a context switch, restore the machine state so that
//this address space can run.
//
      For now, tell the machine where to find the page table.
```

```
void AddrSpace::RestoreState()
   machine->pageTable = pageTable;
   machine->pageTableSize = numPages;
```

```
// addrspace.h
//Data structures to keep track of executing user programs
//(address spaces).
//For now, we don't keep any information about address spaces.
//The user level CPU state is saved and restored in the thread
//executing the user program (see thread.h).
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#ifndef ADDRSPACE H
#define ADDRSPACE_H
#include "copyright.h"
#include "filesys.h"
#define UserStackSize1024 // increase this as necessary!
class AddrSpace {
 public:
   AddrSpace(OpenFile *executable);// Create an address space,
// initializing it with the program
// stored in the file "executable"
   ~AddrSpace();// De-allocate an address space
   void InitRegisters();// Initialize user-level CPU registers,
// before jumping to user code
   void SaveState();// Save/restore address space-specific
   void RestoreState();// info on a context switch
 private:
   TranslationEntry *pageTable;// Assume linear page table translation
// for now!
   unsigned int numPages;// Number of pages in the virtual
// address space
#endif // ADDRSPACE H
```

```
// bitmap.c
//Routines to manage a bitmap -- an array of bits each of which
//can be either on or off. Represented as an array of integers.
// Copyright (c) 1992-1993 The Regents of the University of California. // Return TRUE if the "nth" bit is set.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "bitmap.h"
//-----
// BitMap::BitMap
// Initialize a bitmap with "nitems" bits, so that every bit is clear.
//it can be added somewhere on a list.
//"nitems" is the number of bits in the bitmap.
//-----
BitMap::BitMap(int nitems)
   numBits = nitems;
   numWords = divRoundUp(numBits, BitsInWord);
   map = new unsigned int[numWords];
   for (int i = 0; i < numBits; i++)
      Clear(i);
// BitMap::~BitMap
// De-allocate a bitmap.
//----
BitMap::~BitMap()
   delete map;
,
//-----
// BitMap::Set
// Set the "nth" bit in a bitmap.
//"which" is the number of the bit to be set.
void
BitMap::Mark(int which)
   ASSERT(which >= 0 && which < numBits);
   map[which / BitsInWord] |= 1 << (which % BitsInWord);</pre>
// BitMap::Clear
// Clear the "nth" bit in a bitmap.
//"which" is the number of the bit to be cleared.
BitMap::Clear(int which)
   ASSERT(which >= 0 && which < numBits);
```

```
map[which / BitsInWord] &= ~(1 << (which % BitsInWord));</pre>
.,//-----
// BitMap::Test
//"which" is the number of the bit to be tested.
//-----
bool
BitMap::Test(int which)
   ASSERT(which >= 0 && which < numBits);
   if (map[which / BitsInWord] & (1 << (which % BitsInWord)))</pre>
return TRUE;
   else
return FALSE;
//----
// BitMap::Find
// Return the number of the first bit which is clear.
//As a side effect, set the bit (mark it as in use).
//(In other words, find and allocate a bit.)
//
//If no bits are clear, return -1.
BitMap::Find()
   for (int i = 0; i < numBits; i++)
if (!Test(i)) {
   Mark(i);
   return i;
   return -1;
,
//-----
// BitMap::NumClear
// Return the number of clear bits in the bitmap.
//(In other words, how many bits are unallocated?)
BitMap::NumClear()
   int count = 0;
   for (int i = 0; i < numBits; i++)
if (!Test(i)) count++;
   return count;
// BitMap::Print
// Print the contents of the bitmap, for debugging.
//Could be done in a number of ways, but we just print the #'s of
//all the bits that are set in the bitmap.
```

```
void
BitMap::Print()
   printf("Bitmap set:\n");
   for (int i = 0; i < numBits; i++)
if (Test(i))
   printf("%d, ", i);
   printf("\n");
// These aren't needed until the FILESYS assignment
// BitMap::FetchFromFile
// Initialize the contents of a bitmap from a Nachos file.
//"file" is the place to read the bitmap from
void
BitMap::FetchFrom(OpenFile *file)
   file->ReadAt((char *)map, numWords * sizeof(unsigned), 0);
// BitMap::WriteBack
// Store the contents of a bitmap to a Nachos file.
//"file" is the place to write the bitmap to
BitMap::WriteBack(OpenFile *file)
  file->WriteAt((char *)map, numWords * sizeof(unsigned), 0);
```

```
// bitmap.h
//Data structures defining a bitmap -- an array of bits each of which
//can be either on or off.
//Represented as an array of unsigned integers, on which we do
//modulo arithmetic to find the bit we are interested in.
//The bitmap can be parameterized with with the number of bits being
//managed.
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#ifndef BITMAP H
#define BITMAP H
#include "copyright.h"
#include "utility.h"
#include "openfile.h"
// Definitions helpful for representing a bitmap as an array of
integers
#define BitsInByte 8
#define BitsInWord 32
// The following class defines a "bitmap" -- an array of bits,
// each of which can be independently set, cleared, and tested.
// Most useful for managing the allocation of the elements of an array
// for instance, disk sectors, or main memory pages.
// Each bit represents whether the corresponding sector or page is
// in use or free.
class BitMap {
 public:
   BitMap(int nitems);// Initialize a bitmap, with "nitems" bits
// initially, all bits are cleared.
   ~BitMap();// De-allocate bitmap
   void Mark(int which);
                          // Set the "nth" bit
   void Clear(int which); // Clear the "nth" bit
   bool Test(int which); // Is the "nth" bit set?
   int Find();
                          // Return the # of a clear bit, and as a
side
// effect, set the bit.
// If no bits are clear, return -1.
    int NumClear();// Return the number of clear bits
   void Print();// Print contents of bitmap
   // These aren't needed until FILESYS, when we will need to read and
    // write the bitmap to a file
   void FetchFrom(OpenFile *file); // fetch contents from disk
   void WriteBack(OpenFile *file); // write contents to disk
 private:
   int numBits; // number of bits in the bitmap
   int numWords; // number of words of bitmap storage
// (rounded up if numBits is not a
// multiple of the number of bits in
```

```
// a word)
    unsigned int *map; // bit storage
#endif // BITMAP_H
```

```
// exception.cc
//Entry point into the Nachos kernel from user programs.
//There are two kinds of things that can cause control to
//transfer back to here from user code:
//
//svscall -- The user code explicitly requests to call a procedure
//in the Nachos kernel. Right now, the only function we support is
//"Halt".
//
//exceptions -- The user code does something that the CPU can't handle.
//For instance, accessing memory that doesn't exist, arithmetic errors,
//
//Interrupts (which can also cause control to transfer from user
//code into the Nachos kernel) are handled elsewhere.
//
// For now, this only handles the Halt() system call.
// Everything else core dumps.
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "system.h"
#include "syscall.h"
// ExceptionHandler
// Entry point into the Nachos kernel. Called when a user program
//is executing, and either does a syscall, or generates an addressing
//or arithmetic exception.
//
// For system calls, the following is the calling convention:
// system call code -- r2
//arg1 -- r4
//arg2 -- r5
//arg3 -- r6
//arq4 -- r7
//
//The result of the system call, if any, must be put back into r2.
// And don't forget to increment the pc before returning. (Or else
you'll
// loop making the same system call forever!
//"which" is the kind of exception. The list of possible exceptions
//are in machine.h.
ExceptionHandler(ExceptionType which)
   int type = machine->ReadRegister(2);
   if ((which == SyscallException) && (type == SC_Halt)) {
DEBUG('a', "Shutdown, initiated by user program.\n");
   interrupt->Halt();
```

```
} else {
printf("Unexpected user mode exception %d %d\n", which, type);
ASSERT(FALSE);
```

```
// progtest.cc
//Test routines for demonstrating that Nachos can load
//a user program and execute it.
//
//Also, routines for testing the Console hardware device.
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "system.h"
#include "console.h"
#include "addrspace.h"
#include "synch.h"
//-----
// StartProcess
// Run a user program. Open the executable, load it into
//memory, and jump to it.
void
StartProcess(char *filename)
   OpenFile *executable = fileSystem->Open(filename);
   AddrSpace *space;
   if (executable == NULL) {
printf("Unable to open file %s\n", filename);
return;
   space = new AddrSpace(executable);
   currentThread->space = space;
   delete executable; // close file
   space->InitRegisters();// set the initial register values
   space->RestoreState();// load page table register
   machine->Run();// jump to the user progam
   ASSERT(FALSE);// machine->Run never returns;
// the address space exits
// by doing the syscall "exit"
// Data structures needed for the console test. Threads making
// I/O requests wait on a Semaphore to delay until the I/O completes.
static Console *console;
static Semaphore *readAvail;
static Semaphore *writeDone;
//-----
// ConsoleInterruptHandlers
// Wake up the thread that requested the I/O.
//-----
static void ReadAvail(int arg) { readAvail->V(); }
static void WriteDone(int arg) { writeDone->V(); }
//-----
// ConsoleTest
// Test the console by echoing characters typed at the input onto
//the output. Stop when the user types a 'q'.
void
```

```
/* syscalls.h
 * Nachos system call interface. These are Nachos kernel operations
 * that can be invoked from user programs, by trapping to the kernel
 *via the "syscall" instruction.
 *This file is included by user programs and by the Nachos kernel.
#ifndef SYSCALLS H
#define SYSCALLS H
/* system call codes -- used by the stubs to tell the kernel which
system call
 * is being asked for
 * /
#define SC Halt0
#define SC_Exit1
#define SC_Exec2
#define SC Join3
#define SC_Create4
#define SC_Open5
#define SC Read6
#define SC_Write7
#define SC Close8
#define SC Fork9
#define SC Yield10
#ifndef IN_ASM
/* The system call interface. These are the operations the Nachos
 * kernel needs to support, to be able to run user programs.
 * Each of these is invoked by a user program by simply calling the
 * procedure; an assembly language stub stuffs the system call code
 * into a register, and traps to the kernel. The kernel procedures
 * are then invoked in the Nachos kernel, after appropriate error
checking,
 * from the system call entry point in exception.cc.
/* Stop Nachos, and print out performance stats */
void Halt();
/* Address space control operations: Exit, Exec, and Join */
/* This user program is done (status = 0 means exited normally). */
void Exit(int status);
/* A unique identifier for an executing user program (address space) */
typedef int SpaceId;
/* Run the executable, stored in the Nachos file "name", and return the
 * address space identifier */
SpaceId Exec(char *name);
/* Only return once the the user program "id" has finished.
 * Return the exit status. */
int Join(SpaceId id);
/* File system operations: Create, Open, Read, Write, Close
* These functions are patterned after UNIX -- files represent
 * both files *and* hardware I/O devices.
 * If this assignment is done before doing the file system assignment,
 * note that the Nachos file system has a stub implementation, which
 * will work for the purposes of testing out these routines. */
/* A unique identifier for an open Nachos file. */
typedef int OpenFileId;
/* when an address space starts up, it has two open files, representing
```

```
* keyboard input and display output (in UNIX terms, stdin and
stdout).
* Read and Write can be used directly on these, without first opening
 * the console device. */
#define ConsoleInput0
#define ConsoleOutput1
/* Create a Nachos file, with "name" */
void Create(char *name);
/* Open the Nachos file "name", and return an "OpenFileId" that can
* be used to read and write to the file. */
OpenFileId Open(char *name);
/* Write "size" bytes from "buffer" to the open file. */
void Write(char *buffer, int size, OpenFileId id);
/* Read "size" bytes from the open file into "buffer".
* Return the number of bytes actually read -- if the open file isn't
* long enough, or if it is an I/O device, and there aren't enough
* characters to read, return whatever is available (for I/O devices,
* you should always wait until you can return at least one
character). */
int Read(char *buffer, int size, OpenFileId id);
/* Close the file, we're done reading and writing to it. */
void Close(OpenFileId id);
/* User-level thread operations: Fork and Yield. To allow multiple
* threads to run within a user program. */
/* Fork a thread to run a procedure ("func") in the *same* address
space
* as the current thread. */
void Fork(void (*func)());
/* Yield the CPU to another runnable thread, whether in this address
space
* or not. */
void Yield();
#endif /* IN ASM */
#endif /* SYSCALL_H */
```

```
// console.cc
//Routines to simulate a serial port to a console device.
//A console has input (a keyboard) and output (a display).
//These are each simulated by operations on UNIX files.
//The simulated device is asynchronous.
//so we have to invoke the interrupt handler (after a simulated
//delay), to signal that a byte has arrived and/or that a written
//byte has departed.
//
// DO NOT CHANGE -- part of the machine emulation
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "console.h"
#include "system.h"
// Dummy functions because C++ is weird about pointers to member
functions
static void ConsoleReadPoll(int c)
{ Console *console = (Console *)c; console->CheckCharAvail(); }
static void ConsoleWriteDone(int c)
 Console *console = (Console *)c; console->WriteDone(); }
//-----
// Console::Console
// Initialize the simulation of a hardware console device.
//"readFile" -- UNIX file simulating the keyboard (NULL -> use stdin)
//"writeFile" -- UNIX file simulating the display (NULL -> use stdout)
// "readAvail" is the interrupt handler called when a character arrives ConsoleReadInt);
//from the keyboard
// "writeDone" is the interrupt handler called when a character has
//been output, so that it is ok to request the next char be
//output
Console::Console(char *readFile, char *writeFile, VoidFunctionPtr
readAvail.
VoidFunctionPtr writeDone, int callArg)
   if (readFile == NULL)
readFileNo = 0;// keyboard = stdin
   else
   readFileNo = OpenForReadWrite(readFile, TRUE);// should be read-
only
   if (writeFile == NULL)
writeFileNo = 1;// display = stdout
   else
   writeFileNo = OpenForWrite(writeFile);
   // set up the stuff to emulate asynchronous interrupts
   writeHandler = writeDone;
   readHandler = readAvail;
   handlerArg = callArg;
   putBusy = FALSE;
   incoming = EOF;
   // start polling for incoming packets
```

```
interrupt->Schedule(ConsoleReadPoll, (int)this, ConsoleTime,
ConsoleReadInt);
// Console::~Console
// Clean up console emulation
//-----
Console::~Console()
   if (readFileNo != 0)
Close(readFileNo);
   if (writeFileNo != 1)
Close(writeFileNo);
,
//-----
// Console::CheckCharAvail()
// Periodically called to check if a character is available for
//input from the simulated keyboard (eq, has it been typed?).
//Only read it in if there is buffer space for it (if the previous
//character has been grabbed out of the buffer by the Nachos kernel).
//Invoke the "read" interrupt handler, once the character has been
//put into the buffer.
void
Console::CheckCharAvail()
   char c;
   // schedule the next time to poll for a packet
   interrupt->Schedule(ConsoleReadPoll, (int)this, ConsoleTime,
   // do nothing if character is already buffered, or none to be read
   if ((incoming != EOF) | !PollFile(readFileNo))
return;
   // otherwise, read character and tell user about it
   Read(readFileNo, &c, sizeof(char));
   incoming = c;
   stats->numConsoleCharsRead++;
   (*readHandler)(handlerArg);
//----
// Console::WriteDone()
// Internal routine called when it is time to invoke the interrupt
//handler to tell the Nachos kernel that the output character has
//completed.
//-----
void
Console::WriteDone()
   putBusy = FALSE;
   stats->numConsoleCharsWritten++;
   (*writeHandler)(handlerArg);
.//-----
// Console::GetChar()
// Read a character from the input buffer, if there is any there.
```

```
//Either return the character, or EOF if none buffered.
char
Console::GetChar()
  char ch = incoming;
  incoming = EOF;
  return ch;
   _____
// Console::PutChar()
// Write a character to the simulated display, schedule an interrupt
//to occur in the future, and return.
void
Console::PutChar(char ch)
   ASSERT(putBusy == FALSE);
   WriteFile(writeFileNo, &ch, sizeof(char));
   putBusy = TRUE;
   interrupt->Schedule(ConsoleWriteDone, (int)this, ConsoleTime,
ConsoleWriteInt);
```

```
// console h
//Data structures to simulate the behavior of a terminal
//I/O device. A terminal has two parts -- a keyboard input,
//and a display output, each of which produces/accepts
//characters sequentially.
//The console hardware device is asynchronous. When a character is
//written to the device, the routine returns immediately, and an
//interrupt handler is called later when the I/O completes.
//For reads, an interrupt handler is called when a character arrives.
//The user of the device can specify the routines to be called when
//the read/write interrupts occur. There is a separate interrupt
//for read and write, and the device is "duplex" -- a character
//can be outgoing and incoming at the same time.
//
// DO NOT CHANGE -- part of the machine emulation
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#ifndef CONSOLE H
#define CONSOLE H
#include "copyright.h"
#include "utility.h"
// The following class defines a hardware console device.
// Input and output to the device is simulated by reading
// and writing to UNIX files ("readFile" and "writeFile").
//
// Since the device is asynchronous, the interrupt handler "readAvail"
// is called when a character has arrived, ready to be read in.
// The interrupt handler "writeDone" is called when an output character
// has been "put", so that the next character can be written.
class Console {
 public:
   Console(char *readFile, char *writeFile, VoidFunctionPtr readAvail,
VoidFunctionPtr writeDone, int callArg);
// initialize the hardware console device
    ~Console();// clean up console emulation
// external interface -- Nachos kernel code can call these
    void PutChar(char ch);// Write "ch" to the console display,
// and return immediately. "writeHandler"
// is called when the I/O completes.
    char GetChar();    // Poll the console input. If a char is
// available, return it. Otherwise, return EOF.
    // "readHandler" is called whenever there is
// a char to be gotten
// internal emulation routines -- DO NOT call these.
   void WriteDone(); // internal routines to signal I/O completion
   void CheckCharAvail();
 private:
   int readFileNo; // UNIX file emulating the keyboard
   int writeFileNo; // UNIX file emulating the display
   VoidFunctionPtr writeHandler; // Interrupt handler to call when
// the PutChar I/O completes
```

```
VoidFunctionPtr readHandler; // Interrupt handler to call when // a character arrives from the keyboard int handlerArg;// argument to be passed to the // interrupt handlers bool putBusy; // Is a PutChar operation in progress? // If so, you can't do another one! char incoming; // Contains the character to be read, // if there is one available. // Otherwise contains EOF. }; #endif // CONSOLE_H
```

```
// disk.cc
//Routines to simulate a physical disk device; reading and writing
//to the disk is simulated as reading and writing to a UNIX file.
//See disk.h for details about the behavior of disks (and
//therefore about the behavior of this simulation).
//Disk operations are asynchronous, so we have to invoke an interrupt
//handler when the simulated operation completes.
// DO NOT CHANGE -- part of the machine emulation
//
// Copyright (c) 1992-1993 The Regents of the University of California. //disk.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "disk.h"
#include "system.h"
// We put this at the front of the UNIX file representing the
// disk, to make it less likely we will accidentally treat a useful
file
// as a disk (which would probably trash the file's contents).
#define MagicNumber 0x456789ab
#define MagicSize sizeof(int)
#define DiskSize (MagicSize + (NumSectors * SectorSize))
// dummy procedure because we can't take a pointer of a member function
static void DiskDone(int arg) { ((Disk *)arg)->HandleInterrupt(); }
//-----
// Disk::Disk()
// Initialize a simulated disk. Open the UNIX file (creating it
//if it doesn't exist), and check the magic number to make sure it's
// ok to treat it as Nachos disk storage.
//"name" -- text name of the file simulating the Nachos disk
//"callWhenDone" -- interrupt handler to be called when disk read/write // Disk::ReadRequest/WriteRequest
    request completes
//"callArg" -- argument to pass the interrupt handler
//-----
Disk::Disk(char* name, VoidFunctionPtr callWhenDone, int callArg)
   int magicNum;
   int tmp = 0;
   DEBUG('d', "Initializing the disk, 0x%x 0x%x\n", callWhenDone,
callArg);
   handler = callWhenDone;
   handlerArg = callArg;
   lastSector = 0;
   bufferInit = 0;
   fileno = OpenForReadWrite(name, FALSE);
   if (fileno >= 0) { // file exists, check magic number
Read(fileno, (char *) &magicNum, MagicSize);
ASSERT(magicNum == MagicNumber);
   } else {// file doesn't exist, create it
       fileno = OpenForWrite(name);
magicNum = MagicNumber;
```

```
WriteFile(fileno, (char *) & magicNum, MagicSize); // write magic
number
// need to write at end of file, so that reads will not return EOF
       Lseek(fileno, DiskSize - sizeof(int), 0);
WriteFile(fileno, (char *)&tmp, sizeof(int));
   active = FALSE;
// Disk::~Disk()
// Clean up disk simulation, by closing the UNIX file representing the
Disk::~Disk()
   Close(fileno);
       ______
// Disk::PrintSector()
// Dump the data in a disk read/write request, for debugging.
//-----
static void
PrintSector (bool writing, int sector, char *data)
   int *p = (int *) data;
   if (writing)
       printf("Writing sector: %d\n", sector);
       printf("Reading sector: %d\n", sector);
   for (unsigned int i = 0; i < (SectorSize/sizeof(int)); i++)</pre>
printf("%x ", p[i]);
   printf("\n");
,//-----
// Simulate a request to read/write a single disk sector
    Do the read/write immediately to the UNIX file
    Set up an interrupt handler to be called later,
//
       that will notify the caller when the simulator says
//
       the operation has completed.
//Note that a disk only allows an entire sector to be read/written,
//not part of a sector.
//"sectorNumber" -- the disk sector to read/write
//"data" -- the bytes to be written, the buffer to hold the incoming
Disk::ReadRequest(int sectorNumber, char* data)
   int ticks = ComputeLatency(sectorNumber, FALSE);
   ASSERT(!active);// only one request at a time
   ASSERT((sectorNumber >= 0) && (sectorNumber < NumSectors));
   DEBUG('d', "Reading from sector %d\n", sectorNumber);
```

```
Lseek(fileno, SectorSize * sectorNumber + MagicSize, 0);
   Read(fileno, data, SectorSize);
   if (DebugIsEnabled('d'))
PrintSector(FALSE, sectorNumber, data);
   active = TRUE;
   UpdateLast(sectorNumber);
   stats->numDiskReads++;
   interrupt->Schedule(DiskDone, (int) this, ticks, DiskInt);
void
Disk::WriteRequest(int sectorNumber, char* data)
   int ticks = ComputeLatency(sectorNumber, TRUE);
   ASSERT(!active);
   ASSERT((sectorNumber >= 0) && (sectorNumber < NumSectors));
   DEBUG('d', "Writing to sector %d\n", sectorNumber);
   Lseek(fileno, SectorSize * sectorNumber + MagicSize, 0);
   WriteFile(fileno, data, SectorSize);
   if (DebugIsEnabled('d'))
PrintSector(TRUE, sectorNumber, data);
   active = TRUE;
   UpdateLast(sectorNumber);
   stats->numDiskWrites++;
   interrupt->Schedule(DiskDone, (int) this, ticks, DiskInt);
// Disk::HandleInterrupt()
// Called when it is time to invoke the disk interrupt handler,
//to tell the Nachos kernel that the disk request is done.
void
Disk::HandleInterrupt ()
   active = FALSE;
   (*handler)(handlerArg);
,//-----
// Disk::TimeToSeek()
//Returns how long it will take to position the disk head over the
correct
//track on the disk. Since when we finish seeking, we are likely
//to be in the middle of a sector that is rotating past the head,
//we also return how long until the head is at the next sector
boundary.
//
    Disk seeks at one track per SeekTime ticks (cf. stats.h)
    and rotates at one sector per RotationTime ticks
Disk::TimeToSeek(int newSector, int *rotation)
   int newTrack = newSector / SectorsPerTrack;
   int oldTrack = lastSector / SectorsPerTrack;
```

```
int seek = abs(newTrack - oldTrack) * SeekTime;
// how long will seek take?
   int over = (stats->totalTicks + seek) % RotationTime;
// will we be in the middle of a sector when
// we finish the seek?
   *rotation = 0;
   if (over > 0) // if so, need to round up to next full sector
  *rotation = RotationTime - over;
   return seek;
,
//-----
// Disk::ModuloDiff()
// Return number of sectors of rotational delay between target sector
//"to" and current sector position "from"
//------
Disk::ModuloDiff(int to, int from)
   int toOffset = to % SectorsPerTrack;
   int fromOffset = from % SectorsPerTrack;
   return ((toOffset - fromOffset) + SectorsPerTrack) %
SectorsPerTrack;
í/-----
// Disk::ComputeLatency()
// Return how long will it take to read/write a disk sector, from
//the current position of the disk head.
//
1//
    Latency = seek time + rotational latency + transfer time
    Disk seeks at one track per SeekTime ticks (cf. stats.h)
    and rotates at one sector per RotationTime ticks
//
//
11
    To find the rotational latency, we first must figure out where
the
    disk head will be after the seek (if any). We then figure out
//
//
    how long it will take to rotate completely past newSector after
//that point.
//
    The disk also has a "track buffer"; the disk continuously reads
    the contents of the current disk track into the buffer. This
allows
//
    read requests to the current track to be satisfied more quickly.
    The contents of the track buffer are discarded after every seek
//
to
11
    a new track.
//----
Disk::ComputeLatency(int newSector, bool writing)
   int rotation;
   int seek = TimeToSeek(newSector, &rotation);
   int timeAfter = stats->totalTicks + seek + rotation;
#ifndef NOTRACKBUF// turn this on if you don't want the track buffer
stuff
   // check if track buffer applies
   if ((writing == FALSE) && (seek == 0)
```

```
&& (((timeAfter - bufferInit) / RotationTime)
    > ModuloDiff(newSector, bufferInit / RotationTime))) {
       DEBUG('d', "Request latency = %d\n", RotationTime);
return RotationTime; // time to transfer sector from the track buffer
#endif
   rotation += ModuloDiff(newSector, timeAfter / RotationTime) *
RotationTime;
   DEBUG('d', "Request latency = %d\n", seek + rotation +
RotationTime);
   return(seek + rotation + RotationTime);
   ______
// Disk::UpdateLast
    Keep track of the most recently requested sector. So we can know
//what is in the track buffer.
void
Disk::UpdateLast(int newSector)
   int rotate;
   int seek = TimeToSeek(newSector, &rotate);
   if (seek != 0)
bufferInit = stats->totalTicks + seek + rotate;
   lastSector = newSector;
   DEBUG('d', "Updating last sector = %d, %d\n", lastSector,
bufferInit);
```

```
// disk.h
//Data structures to emulate a physical disk. A physical disk
//can accept (one at a time) requests to read/write a disk sector;
//when the request is satisfied, the CPU gets an interrupt, and
//the next request can be sent to the disk.
//Disk contents are preserved across machine crashes, but if
//a file system operation (eq. create a file) is in progress when the
//system shuts down, the file system may be corrupted.
//
// DO NOT CHANGE -- part of the machine emulation
#ifndef DISK H
#define DISK H
#include "copyright.h"
#include "utility.h"
// The following class defines a physical disk I/O device. The disk
// has a single surface, split up into "tracks", and each track split
// up into "sectors" (the same number of sectors on each track, and
each
// sector has the same number of bytes of storage).
//
// Addressing is by sector number -- each sector on the disk is given
// a unique number: track * SectorsPerTrack + offset within a track.
//
// As with other I/O devices, the raw physical disk is an asynchronous
device --
// requests to read or write portions of the disk return immediately,
// and an interrupt is invoked later to signal that the operation
completed.
//
// The physical disk is in fact simulated via operations on a UNIX
file.
//
// To make life a little more realistic, the simulated time for
// each operation reflects a "track buffer" -- RAM to store the
contents
// of the current track as the disk head passes by. The idea is that
the
// disk always transfers to the track buffer, in case that data is
requested
// later on. This has the benefit of eliminating the need for
// "skip-sector" scheduling -- a read request which comes in shortly
after
// the head has passed the beginning of the sector can be satisfied
// quickly, because its contents are in the track buffer. Most
// disks these days now come with a track buffer.
//
// The track buffer simulation can be disabled by compiling with -
DNOTRACKBUF
#define SectorSize 128// number of bytes per disk sector
#define SectorsPerTrack 32// number of sectors per disk track
#define NumTracks 32// number of tracks per disk
#define NumSectors (SectorsPerTrack * NumTracks)
// total # of sectors per disk
class Disk {
```

```
public:
    Disk(char* name, VoidFunctionPtr callWhenDone, int callArg);
    // Create a simulated disk.
// Invoke (*callWhenDone)(callArg)
// every time a request completes.
    ~Disk();// Deallocate the disk.
    void ReadRequest(int sectorNumber, char* data);
    // Read/write an single disk sector.
// These routines send a request to
    // the disk and return immediately.
    // Only one request allowed at a time!
    void WriteRequest(int sectorNumber, char* data);
    void HandleInterrupt();// Interrupt handler, invoked when
// disk request finishes.
    int ComputeLatency(int newSector, bool writing);
    // Return how long a request to
// newSector will take:
// (seek + rotational delay + transfer)
 private:
    int fileno; // UNIX file number for simulated disk
    VoidFunctionPtr handler; // Interrupt handler, to be invoked
// when any disk request finishes
    int handlerArg; // Argument to interrupt handler
                     // Is a disk operation in progress?
    bool active;
    int lastSector;// The previous disk request
    int bufferInit; // When the track buffer started
// being loaded
    int TimeToSeek(int newSector, int *rotate); // time to get to the
new track
    int ModuloDiff(int to, int from);
                                             // # sectors between to
and from
    void UpdateLast(int newSector);
#endif // DISK_H
```

```
// machine.cc
//Routines for simulating the execution of user programs.
// DO NOT CHANGE -- part of the machine emulation
#include "copyright.h"
#include "machine.h"
#include "system.h"
// Textual names of the exceptions that can be generated by user
// execution, for debugging.
static char* exceptionNames[] = { "no exception", "syscall",
"page fault/no TLB entry", "page read only",
"bus error", "address error", "overflow",
"illegal instruction" };
//----
// CheckEndian
// Check to be sure that the host really uses the format it says it
//does, for storing the bytes of an integer. Stop on error.
//-----
static
void CheckEndian()
   union checkit {
      char charword[4];
      unsigned int intword;
   } check;
   check.charword[0] = 1;
   check.charword[1] = 2;
   check.charword[2] = 3;
   check.charword[3] = 4;
#ifdef HOST IS BIG ENDIAN
   ASSERT (check.intword == 0x01020304);
#else
   ASSERT (check.intword == 0 \times 04030201);
#endif
  ______
// Machine::Machine
// Initialize the simulation of user program execution.
//"debug" -- if TRUE, drop into the debugger after each user
instruction
//is executed.
//----
Machine::Machine(bool debug)
   int i;
   for (i = 0; i < NumTotalRegs; i++)
      registers[i] = 0;
   mainMemory = new char[MemorySize];
   for (i = 0; i < MemorySize; i++)</pre>
     mainMemory[i] = 0;
#ifdef USE_TLB
   tlb = new TranslationEntry[TLBSize];
   for (i = 0; i < TLBSize; i++)
tlb[i].valid = FALSE;
```

```
pageTable = NULL;
#else// use linear page table
   tlb = NULL;
   pageTable = NULL;
#endif
   singleStep = debug;
   CheckEndian();
,
//-----
// Machine::~Machine
// De-allocate the data structures used to simulate user program
//-----
Machine::~Machine()
   delete [] mainMemory;
   if (tlb != NULL)
      delete [] tlb;
// Machine::RaiseException
// Transfer control to the Nachos kernel from user mode, because
//the user program either invoked a system call, or some exception
//occured (such as the address translation failed).
//"which" -- the cause of the kernel trap
//"badVaddr" -- the virtual address causing the trap, if appropriate
//-----
biov
Machine::RaiseException(ExceptionType which, int badVAddr)
   DEBUG('m', "Exception: %s\n", exceptionNames[which]);
// ASSERT(interrupt->getStatus() == UserMode);
   registers[BadVAddrReg] = badVAddr;
   DelayedLoad(0, 0);// finish anything in progress
   interrupt->setStatus(SystemMode);
   ExceptionHandler(which);// interrupts are enabled at this point
   interrupt->setStatus(UserMode);
//-----
// Machine::Debugger
// Primitive debugger for user programs. Note that we can't use
//qdb to debug user programs, since qdb doesn't run on top of Nachos.
//It could, but you'd have to implement *a lot* more system calls
//to get it to work!
//So just allow single-stepping, and printing the contents of memory.
void Machine::Debugger()
   char *buf = new char[80];
   int num;
   interrupt->DumpState();
   DumpState();
   printf("%d> ", stats->totalTicks);
```

```
fflush(stdout);
   fgets(buf, 80, stdin);
   if (sscanf(buf, "%d", &num) == 1)
runUntilTime = num;
   else {
runUntilTime = 0;
switch (*buf) {
 case '\n':
   break;
 case 'c':
   singleStep = FALSE;
   break;
 case '?':
   printf("Machine commands:\n");
             <return> execute one instruction\n");
   printf("
   printf("
              <number> run until the given timer tick\n");
   printf("
              c run until completion\n");
              ? print help message\n");
   printf("
   break;
   delete [] buf;
//-----
// Machine::DumpState
// Print the user program's CPU state. We might print the contents
//of memory, but that seemed like overkill.
//----
biov
Machine::DumpState()
   int i;
   printf("Machine registers:\n");
   for (i = 0; i < NumGPRegs; i++)
switch (i) {
 case StackReq:
   printf("\tSP(%d):\t0x%x%s", i, registers[i],
  ((i % 4) == 3) ? "\n" : "");
   break;
 case RetAddrReq:
   printf("\tRA(%d):\t0x%x%s", i, registers[i],
  ((i % 4) == 3) ? "\n" : "");
   break;
 default:
   printf("\t%d:\t0x%x%s", i, registers[i],
  ((i % 4) == 3) ? "\n" : "");
   break;
   printf("\tHi:\t0x%x", registers[HiReq]);
```

```
printf("\tLo:\t0x%x\n", registers[LoReq]);
   printf("\tPC:\t0x%x", registers[PCReq]);
   printf("\tNextPC:\t0x%x", registers[NextPCReq]);
   printf("\tPrevPC:\t0x%x\n", registers[PrevPCReq]);
   printf("\tLoad:\t0x%x", registers[LoadReg]);
   printf("\tLoadV:\t0x%x\n", registers[LoadValueReg]);
   printf("\n");
// Machine::ReadRegister/WriteRegister
// Fetch or write the contents of a user program register.
//-----
int Machine::ReadRegister(int num)
ASSERT((num >= 0) && (num < NumTotalRegs));
return registers[num];
void Machine::WriteRegister(int num, int value)
ASSERT((num >= 0) && (num < NumTotalRegs));
// DEBUG('m', "WriteRegister %d, value %d\n", num, value);
registers[num] = value;
```

```
// machine.h
//Data structures for simulating the execution of user programs
//running on top of Nachos.
//User programs are loaded into "mainMemory"; to Nachos,
//this looks just like an array of bytes. Of course, the Nachos
//kernel is in memory too -- but as in most machines these days,
//the kernel is loaded into a separate memory region from user
//programs, and accesses to kernel memory are not translated or paged.
//In Nachos, user programs are executed one instruction at a time,
//by the simulator. Each memory reference is translated, checked
//for errors, etc.
#ifndef MACHINE H
#define MACHINE H
#include "copyright.h"
#include "utility.h"
#include "translate.h"
#include "disk.h"
// Definitions related to the size, and format of user memory
#define PageSize SectorSize // set the page size equal to
// the disk sector size, for simplicity
#define NumPhysPages
                        32
#define MemorySize (NumPhysPages * PageSize)
#define TLBSize4// if there is a TLB, make it small
enum ExceptionType { NoException,
                                            // Everything ok!
     SyscallException,
                            // A program executed a system call.
     PageFaultException,
                            // No valid translation found
     ReadOnlyException,
                            // Write attempted to page marked "read-
only"
                            // Translation resulted in an
     BusErrorException.
// invalid physical address
     AddressErrorException, // Unaligned reference or one that
// was beyond the end of the address space
     OverflowException,
                           // Integer overflow in add or sub.
     IllegalInstrException, // Unimplemented or reserved instr.
     NumExceptionTypes
// User program CPU state. The full set of MIPS registers, plus a few |// Routines internal to the machine simulation -- DO NOT call these
// more because we need to be able to start/stop a user program between
// any two instructions (thus we need to keep track of things like load
// delay slots, etc.)
#define StackReg29// User's stack pointer
#define RetAddrReg31// Holds return address for procedure calls
#define NumGPRegs32// 32 general purpose registers on MIPS
#define HiReq32// Double register to hold multiply result
#define LoReq33
#define PCReq34// Current program counter
#define NextPCReq35// Next program counter (for branch delay)
#define PrevPCReq36// Previous program counter (for debugging)
#define LoadReg37// The register target of a delayed load.
#define LoadValueReq 38// The value to be loaded by a delayed load.
#define BadVAddrReg39// The failing virtual address on an exception
#define NumTotalRegs 40
// The following class defines an instruction, represented in both
// undecoded binary form decoded to identify operation to do
//registers to act on any immediate operand value
class Instruction {
```

```
public:
    void Decode();// decode the binary representation of the
instruction
    unsigned int value; // binary representation of the instruction
    char opCode;
                     // Type of instruction. This is NOT the same as
the
         // opcode field from the instruction: see defs in mips.h
    char rs, rt, rd; // Three registers from instruction.
    int extra;
                     // Immediate or target or shamt field or offset.
                     // Immediates are sign-extended.
// The following class defines the simulated host workstation
hardware, as
// seen by user programs -- the CPU registers, main memory, etc.
// User programs shouldn't be able to tell that they are running on
our
// simulator or on the real hardware, except
//we don't support floating point instructions
//the system call interface to Nachos is not the same as UNIX
// (10 system calls in Nachos vs. 200 in UNIX!)
// If we were to implement more of the UNIX system calls, we ought to
// able to run Nachos on top of Nachos!
// The procedures in this class are defined in machine.cc, mipssim.cc,
and
// translate.cc.
class Machine {
  public:
    Machine(bool debug);// Initialize the simulation of the hardware
// for running user programs
    ~Machine();// De-allocate the data structures
// Routines callable by the Nachos kernel
    void Run(); // Run a user program
    int ReadRegister(int num);// read the contents of a CPU register
    void WriteRegister(int num, int value);
// store a value into a CPU register
    void OneInstruction(Instruction *instr);
    // Run one instruction of a user program.
    void DelayedLoad(int nextReg, int nextVal);
// Do a pending delayed load (modifying a reg)
    bool ReadMem(int addr, int size, int* value);
    bool WriteMem(int addr, int size, int value);
    // Read or write 1, 2, or 4 bytes of virtual
// memory (at addr). Return FALSE if a
// correct translation couldn't be found.
    ExceptionType Translate(int virtAddr, int* physAddr, int size,bool
writing);
    // Translate an address, and check for
// alignment. Set the use and dirty bits in
// the translation entry appropriately,
    // and return an exception code if the
// translation couldn't be completed.
    void RaiseException(ExceptionType which, int badVAddr);
```

```
// Trap to the Nachos kernel, because of a
// system call or other exception.
   void Debugger();// invoke the user program debugger
   void DumpState();// print the user CPU and memory state
// Data structures -- all of these are accessible to Nachos kernel
// "public" for convenience.
// Note that *all* communication between the user program and the
// are in terms of these data structures.
    char *mainMemory; // physical memory to store user program,
// code and data, while executing
    int registers[NumTotalRegs]; // CPU registers, for executing user
programs
// NOTE: the hardware translation of virtual addresses in the user
program
// to physical addresses (relative to the beginning of "mainMemory")
// can be controlled by one of:
//a traditional linear page table
// a software-loaded translation lookaside buffer (tlb) -- a cache of
// mappings of virtual page #'s to physical page #'s
// If "tlb" is NULL, the linear page table is used
// If "tlb" is non-NULL, the Nachos kernel is responsible for managing
//the contents of the TLB. But the kernel can use any data structure
//it wants (eq, segmented paging) for handling TLB cache misses.
// For simplicity, both the page table pointer and the TLB pointer are
// public. However, while there can be multiple page tables (one per
address
// space, stored in memory), there is only one TLB (implemented in
hardware).
// Thus the TLB pointer should be considered as *read-only*, although
// the contents of the TLB are free to be modified by the kernel
software.
    TranslationEntry *tlb;// this pointer should be considered
// "read-only" to Nachos kernel code
   TranslationEntry *pageTable;
   unsigned int pageTableSize;
 private:
   bool singleStep;// drop back into the debugger after each
// simulated instruction
    int runUntilTime; // drop back into the debugger when simulated
// time reaches this value
};
extern void ExceptionHandler(ExceptionType which);
// Entry point into Nachos for handling
// user system calls and exceptions
// Defined in exception.cc
// Routines for converting Words and Short Words to and from the
// simulated machine's format of little endian. If the host machine
// is little endian (DEC and Intel), these end up being NOPs.
// What is stored in each format:
//host byte ordering:
    kernel data structures
     user registers
//simulated machine byte ordering:
     contents of main memory
```

```
unsigned int WordToHost(unsigned int word);
unsigned short ShortToHost(unsigned short shortword);
unsigned int WordToMachine(unsigned int word);
unsigned short ShortToMachine(unsigned short shortword);
#endif // MACHINE H
```

```
// timer cc
//Routines to emulate a hardware timer device.
//
       A hardware timer generates a CPU interrupt every X
milliseconds.
       This means it can be used for implementing time-slicing.
//
//
       We emulate a hardware timer by scheduling an interrupt to occur
//
       every time stats->totalTicks has increased by TimerTicks.
//
       In order to introduce some randomness into time-slicing, if
//
"doRandom"
//
       is set, then the interrupt is comes after a random number of
ticks.
//
//Remember -- nothing in here is part of Nachos. It is just
//an emulation for the hardware that Nachos is running on top of.
// DO NOT CHANGE -- part of the machine emulation
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "timer.h"
#include "system.h"
// dummy function because C++ does not allow pointers to member
functions
static void TimerHandler(int arg)
{ Timer *p = (Timer *)arg; p->TimerExpired(); }
// Timer::Timer
       Initialize a hardware timer device. Save the place to call
//on each interrupt, and then arrange for the timer to start
//generating interrupts.
//
       "timerHandler" is the interrupt handler for the timer device.
//
//It is called with interrupts disabled every time the
//the timer expires.
//
       "callArg" is the parameter to be passed to the interrupt
handler.
//
       "doRandom" -- if true, arrange for the interrupts to occur
//at random, instead of fixed, intervals.
//-----
Timer::Timer(VoidFunctionPtr timerHandler, int callArg, bool doRandom)
   randomize = doRandom;
   handler = timerHandler;
   arg = callArg;
   // schedule the first interrupt from the timer device
   interrupt->Schedule(TimerHandler, (int) this,
TimeOfNextInterrupt(),
TimerInt);
```

```
// Timer::TimerExpired
//
       Routine to simulate the interrupt generated by the hardware
//timer device. Schedule the next interrupt, and invoke the
//interrupt handler.
biov
Timer::TimerExpired()
   // schedule the next timer device interrupt
   interrupt->Schedule(TimerHandler, (int) this,
TimeOfNextInterrupt(),
TimerInt);
    // invoke the Nachos interrupt handler for this device
   (*handler)(arg);
//-----
// Timer::TimeOfNextInterrupt
       Return when the hardware timer device will next cause an
interrupt.
//If randomize is turned on, make it a (pseudo-)random delay.
Timer::TimeOfNextInterrupt()
   if (randomize)
return 1 + (Random() % (TimerTicks * 2));
   else
return TimerTicks;
```

```
// timer h
//Data structures to emulate a hardware timer.
//A hardware timer generates a CPU interrupt every X milliseconds.
//This means it can be used for implementing time-slicing, or for
//having a thread go to sleep for a specific period of time.
//We emulate a hardware timer by scheduling an interrupt to occur
//every time stats->totalTicks has increased by TimerTicks.
//In order to introduce some randomness into time-slicing, if
//is set, then the interrupt comes after a random number of ticks.
//
   DO NOT CHANGE -- part of the machine emulation
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#ifndef TIMER H
#define TIMER H
#include "copyright.h"
#include "utility.h"
// The following class defines a hardware timer.
class Timer {
 public:
   Timer(VoidFunctionPtr timerHandler, int callArg, bool doRandom);
// Initialize the timer, to call the interrupt
// handler "timerHandler" every time slice.
    ~Timer() {}
// Internal routines to the timer emulation -- DO NOT call these
   void TimerExpired();// called internally when the hardware
// timer generates an interrupt
   int TimeOfNextInterrupt(); // figure out when the timer will
generate
// its next interrupt
 private:
   bool randomize;// set if we need to use a random timeout delay
   VoidFunctionPtr handler;// timer interrupt handler
   int arg; // argument to pass to interrupt handler
#endif // TIMER_H
```

```
// translate.cc
//Routines to translate virtual addresses to physical addresses.
//Software sets up a table of legal translations. We look up
//in the table on every memory reference to find the true physical
//memory location.
// Two types of translation are supported here.
//
//Linear page table -- the virtual page # is used as an index
//into the table, to find the physical page #.
//Translation lookaside buffer -- associative lookup in the table
//to find an entry with the same virtual page #. If found,
//this entry is used for the translation.
//If not, it traps to software with an exception.
//
//In practice, the TLB is much smaller than the amount of physical
//memory (16 entries is common on a machine that has 1000's of
//pages). Thus, there must also be a backup translation scheme
//(such as page tables), but the hardware doesn't need to know
//anything at all about that.
//Note that the contents of the TLB are specific to an address space.
//If the address space changes, so does the contents of the TLB!
// DO NOT CHANGE -- part of the machine emulation
//
// Copyright (c) 1992-1993 The Regents of the University of California. Machine::ReadMem(int addr, int size, int *value)
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#include "copyright.h"
#include "machine.h"
#include "addrspace.h"
#include "system.h"
// Routines for converting Words and Short Words to and from the
// simulated machine's format of little endian. These end up
// being NOPs when the host machine is also little endian (DEC and
Intel).
unsigned int
WordToHost(unsigned int word) {
#ifdef HOST IS BIG ENDIAN
 register unsigned long result;
 result = (word >> 24) & 0 \times 00000000ff;
 result |= (word >> 8) & 0x0000ff00;
 result |= (word << 8) & 0x00ff0000;
 result |= (word << 24) & 0xff000000;
 return result;
#else
return word;
#endif /* HOST_IS_BIG_ENDIAN */
unsigned short
ShortToHost(unsigned short shortword) {
#ifdef HOST IS BIG ENDIAN
register unsigned short result;
```

```
result = (shortword << 8) & 0xff00;
 result |= (shortword >> 8) & 0x00ff;
 return result;
#else
 return shortword;
#endif /* HOST IS BIG ENDIAN */
unsigned int
WordToMachine(unsigned int word) { return WordToHost(word); }
unsigned short
ShortToMachine(unsigned short shortword) { return
ShortToHost(shortword); }
//-----
// Machine::ReadMem
        Read "size" (1, 2, or 4) bytes of virtual memory at "addr"
into
//the location pointed to by "value".
// Returns FALSE if the translation step from virtual to physical
memory
// failed.
11
//"addr" -- the virtual address to read from
//"size" -- the number of bytes to read (1, 2, or 4)
//"value" -- the place to write the result
bool
   int data;
    ExceptionType exception;
    int physicalAddress;
    DEBUG('a', "Reading VA 0x%x, size %d\n", addr, size);
    exception = Translate(addr, &physicalAddress, size, FALSE);
    if (exception != NoException) {
machine->RaiseException(exception, addr);
return FALSE;
    switch (size) {
      case 1:
data = machine->mainMemory[physicalAddress];
*value = data;
break;
      case 2:
data = *(unsigned short *) &machine->mainMemory[physicalAddress];
*value = ShortToHost(data);
break;
      case 4:
data = *(unsigned int *) &machine->mainMemory[physicalAddress];
*value = WordToHost(data);
break;
      default: ASSERT(FALSE);
```

```
DEBUG('a', "\tvalue read = %8.8x\n", *value);
   return (TRUE);
// Machine::WriteMem
       Write "size" (1, 2, or 4) bytes of the contents of "value" into ExceptionType
//virtual memory at location "addr".
//
    Returns FALSE if the translation step from virtual to physical
memory
//
    failed.
//
//"addr" -- the virtual address to write to
//"size" -- the number of bytes to be written (1, 2, or 4)
//"value" -- the data to be written
Machine::WriteMem(int addr, int size, int value)
   ExceptionType exception;
   int physicalAddress;
   DEBUG('a', "Writing VA 0x%x, size %d, value 0x%x\n", addr, size,
value);
   exception = Translate(addr, &physicalAddress, size, TRUE);
   if (exception != NoException) {
machine->RaiseException(exception, addr);
return FALSE;
   switch (size) {
     case 1:
machine->mainMemory[physicalAddress] = (unsigned char) (value & 0xff); | if (vpn >= pageTableSize) {
break;
     case 2:
*(unsigned short *) &machine->mainMemory[physicalAddress]
= ShortToMachine((unsigned short) (value & 0xffff));
break;
     case 4:
*(unsigned int *) &machine->mainMemory[physicalAddress]
= WordToMachine((unsigned int) value);
break;
     default: ASSERT(FALSE);
   return TRUE;
·//-----
// Machine::Translate
// Translate a virtual address into a physical address, using
//either a page table or a TLB. Check for alignment and all sorts
//of other errors, and if everything is ok, set the use/dirty bits in
//the translation table entry, and store the translated physical
//address in "physAddr". If there was an error, returns the type
//of the exception.
```

```
//"virtAddr" -- the virtual address to translate
//"physAddr" -- the place to store the physical address
//"size" -- the amount of memory being read or written
// "writing" -- if TRUE, check the "read-only" bit in the TLB
//-----
Machine::Translate(int virtAddr, int* physAddr, int size, bool
writing)
    int i;
   unsigned int vpn, offset;
    TranslationEntry *entry;
   unsigned int pageFrame;
   DEBUG('a', "\tTranslate 0x%x, %s: ", virtAddr, writing ? "write" :
"read");
// check for alignment errors
    if (((size == 4) && (virtAddr & 0x3)) | | ((size == 2) && (virtAddr
& 0x1))){
DEBUG('a', "alignment problem at %d, size %d!\n", virtAddr, size);
return AddressErrorException;
    // we must have either a TLB or a page table, but not both!
   ASSERT(tlb == NULL | pageTable == NULL);
   ASSERT(tlb != NULL | pageTable != NULL);
// calculate the virtual page number, and offset within the page,
// from the virtual address
    vpn = (unsigned) virtAddr / PageSize;
    offset = (unsigned) virtAddr % PageSize;
    if (tlb == NULL) {// => page table => vpn is index into table
    DEBUG('a', "virtual page # %d too large for page table size
%d!\n",
virtAddr, pageTableSize);
    return AddressErrorException;
} else if (!pageTable[vpn].valid) {
    DEBUG('a', "virtual page # %d too large for page table size
%d!\n",
virtAddr, pageTableSize);
   return PageFaultException;
entry = &pageTable[vpn];
   } else {
        for (entry = NULL, i = 0; i < TLBSize; i++)</pre>
       if (tlb[i].valid && (tlb[i].virtualPage == vpn)) {
entry = &tlb[i];// FOUND!
break;
if (entry == NULL) {// not found
       DEBUG('a', "*** no valid TLB entry found for this virtual
       return PageFaultException; // really, this is a TLB fault,
// the page may be in memory,
// but not in the TLB
```

```
if (entry->readOnly && writing) {// trying to write to a read-only
DEBUG('a', "%d mapped read-only at %d in TLB!\n", virtAddr, i);
return ReadOnlyException;
   pageFrame = entry->physicalPage;
   // if the pageFrame is too big, there is something really wrong!
   // An invalid translation was loaded into the page table or TLB.
   if (pageFrame >= NumPhysPages) {
DEBUG('a', "*** frame %d > %d!\n", pageFrame, NumPhysPages);
return BusErrorException;
   entry->use = TRUE;// set the use, dirty bits
    if (writing)
entry->dirty = TRUE;
    *physAddr = pageFrame * PageSize + offset;
   ASSERT((*physAddr >= 0) && ((*physAddr + size) <= MemorySize));
   DEBUG('a', "phys addr = 0x%x\n", *physAddr);
    return NoException;
```

```
// translate.h
//Data structures for managing the translation from
//virtual page # -> physical page #, used for managing
//physical memory on behalf of user programs.
//
//The data structures in this file are "dual-use" - they
//serve both as a page table entry, and as an entry in
//a software-managed translation lookaside buffer (TLB).
//Either way, each entry is of the form:
//<virtual page #, physical page #>.
//
// DO NOT CHANGE -- part of the machine emulation
//
// Copyright (c) 1992-1993 The Regents of the University of California.
// All rights reserved. See copyright.h for copyright notice and
limitation
// of liability and disclaimer of warranty provisions.
#ifndef TLB H
#define TLB H
#include "copyright.h"
#include "utility.h"
// The following class defines an entry in a translation table --
// in a page table or a TLB. Each entry defines a mapping from one
// virtual page to one physical page.
// In addition, there are some extra bits for access control (valid and
// read-only) and some bits for usage information (use and dirty).
class TranslationEntry {
 public:
   int virtualPage; // The page number in virtual memory.
   int physicalPage; // The page number in real memory (relative to
the
// start of "mainMemory"
                        // If this bit is set, the translation is
   bool valid;
ignored.
// (In other words, the entry hasn't been initialized.)
   bool readOnly; // If this bit is set, the user program is not
allowed
// to modify the contents of the page.
   bool use;
                        // This bit is set by the hardware every time
the
// page is referenced or modified.
   bool dirty;
                        // This bit is set by the hardware every time
the
// page is modified.
};
#endif
```

```
#include "syscall.h"
int
main()
    SpaceId newProc;
    OpenFileId input = ConsoleInput;
    OpenFileId output = ConsoleOutput;
    char prompt[2], ch, buffer[60];
    int i;
    prompt[0] = '-';
    prompt[1] = '-';
    while(1)
Write(prompt, 2, output);
i = 0;
do {
    Read(&buffer[i], 1, input);
} while( buffer[i++] != '\n' );
buffer[--i] = '\setminus 0';
if( i > 0 ) {
newProc = Exec(buffer);
Join(newProc);
```

```
/* Start.s
 *Assembly language assist for user programs running on top of Nachos.
 *Since we don't want to pull in the entire C library, we define
 *what we need for a user program here, namely Start and the system
 *calls.
 * /
#define IN ASM
#include "syscall.h"
       .text
        .aliqn 2
  __start
 *Initialize running a C program, by calling "main".
 * NOTE: This has to be first, so that it gets loaded at location 0.
 *The Nachos kernel always starts a program by jumping to location 0.
.globl __start
.ent start
start:
ialmain
move$4,$0
jalExit /* if we return from main, exit(0) */
.end start
/* -----
 * System call stubs:
 *Assembly language assist to make system calls to the Nachos kernel.
 *There is one stub per system call, that places the code for the
 *system call into register r2, and leaves the arguments to the
 *system call alone (in other words, arg1 is in r4, arg2 is
 *in r5, arg3 is in r6, arg4 is in r7)
 * The return value is in r2. This follows the standard C calling
 * convention on the MIPS.
.qlobl Halt
.entHalt
Halt:
addiu $2,$0,SC Halt
syscall
i$31
.end Halt
.qlobl Exit
.entExit
Exit:
addiu $2,$0,SC_Exit
svscall
j$31
.end Exit
.globl Exec
.entExec
Exec:
addiu $2,$0,SC_Exec
```

syscall i\$31 .end Exec .alobl Join .ent.Toin Join: addiu \$2,\$0,SC Join svscall i\$31 end Join .globl Create .entCreate Create: addiu \$2,\$0,SC Create syscall j\$31 .end Create .globl Open .entOpen Open: addiu \$2,\$0,SC Open syscall i\$31 .end Open .qlobl Read .entRead Read: addiu \$2,\$0,SC Read syscall j\$31 .end Read .globl Write .entWrite Write: addiu \$2,\$0,SC_Write syscall j\$31 .end Write .qlobl Close .entClose Close: addiu \$2,\$0,SC Close syscall i\$31 .end Close .qlobl Fork .entFork Fork: addiu \$2,\$0,SC_Fork svscall i\$31 .end Fork .globl Yield .entYield Yield: addiu \$2,\$0,SC_Yield

P#1:1-threads P#1 : ** list.cc P#3 : ** list.h P#4 : ** main.cc P#6: ** scheduler.cc P#8: ** scheduler.h P#9 : ** switch.c P#11: ** switch.h P#13: ** synch.cc P#14: ** synch.h P#16: ** synchlist.cc P#17: ** synchlist.h P#18: ** system.cc P#20 : ** system.h P#21 : ** thread.cc P#24: ** thread.h P#26: ** threadtest.cc P#27 : ** utility.cc P#28: ** utility.h P#29: 2-filesys P#29: ** directory.cc P#31: ** directory.h P#32: ** filehdr.cc P#34: ** filehdr.h P#35: ** filesys.cc P#38: ** filesys.h P#39 : ** fstest.cc P#41: ** openfile.cc P#43: ** openfile.h P#44: ** synchdisk.cc P#45: ** synchdisk.h P#46: 3-userprog P#46: ** addrspace.cc P#48: ** addrspace.h P#49 : ** bitmap.cc

P#51: ** bitmap.h P#52: ** exception.cc P#53: ** progtest.cc P#54: ** syscall.h P#55: 4-machine P#55 : ** console.cc P#57: ** console.h P#58 : ** disk.cc P#61 : ** disk.h P#62: ** machine.cc P#64: ** machine.h P#66 : ** timer.cc P#67: ** timer.h P#68: ** translate.cc P#71: ** translate.h P#72 : 5-test P#72: ** shell.c P#73 : ** start.c