#### 操作系统研讨课

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# Lecture 2 Non-Preemptive Kernel

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#### Schedule

- Project 1 due
- Project 2 assignment



## Project 1 Due

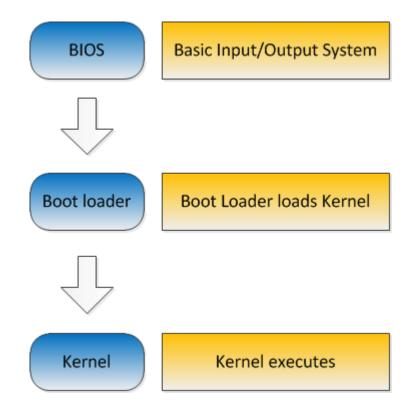
- P1 due
  - We test Tasks 3 and 4 for P1 due
  - Please compile your code, running the code on your board, and show the results to TA
  - Answer any questions we may ask

## Project 1 Due

- P1 submission
  - Submit a compressed package named as "StudentNo.-YourName-P1"
  - Please includes
    - Source code
    - README to simply describe your code, e.g. which file is your work
    - Design document covering the questions in the design document template
  - Do not forget to submit before 23:59 tonight



Booting procedure





- Requirements
  - Write a non-preemptive kernel
    - Start a set of processes/threads
    - Perform context switches between processes and threads
    - Measure the cost of context switch
    - Implement basic mutual exclusion
    - Package the binary of kernel and processes/threads into machine image



- Assumption about multi-tasks
  - A set of tasks
    - Program codes in processX.c and thX.c
    - Test defined in tasks.c
    - Fixed number of tasks for each test
    - Allocate per-task state statically in kernel.c
  - Non-preemptive tasks
    - Run until the task voluntarily yield or exit



- Process Control Block (PCB, TCB)
  - A data structure in OS kernel containing the information to manage a particular process
  - Central role in process management
  - Kept in an area of memory protected from normal user access

- Process Control Block (PCB, TCB)
  - Process identification data
    - process\_type
  - Process state data
    - Status of a process when it is suspended
    - Contents of registers, stack pointers etc.
  - Process control data
    - Process scheduling state
      - process\_state



- Start a task How to describe a task?
  - Task type
    - Process
      - Compiled separately from the kernel
      - Share the same address with the kernel
      - Only setup one stack
      - Still use faked SYSCALL to call kernel functions
    - Kernel thread
      - Share the same address with the kernel
  - Task entry point
    - Function addresses



- Start a task
  - How to associate each task with a PCB?
  - Initialize PCB
    - Which registers should be set?
    - Where is the task located?
    - How to setup stack? Stack size?
    - Where is the PCB located?



- Start a task
  - Possible memory layout
  - Decide your ownSTACK\_MIN and STACK\_MAX

kseg1 Unmapped Uncachable 0xE000 0000

STACK\_MAX

STACK\_MIN

0xA083 0000 (process3 addr.)

0xA082 0000 (process2 addr.)

0xA081 0000 (process1 addr.)

0xA080 026C (kernel entry addr.)

0xA080 0200 (kernel loaded addr.)

0xA080 0030 (bootblock exec addr.)

0xA080 0000 (bootblock loaded addr.)

0xA000 0000



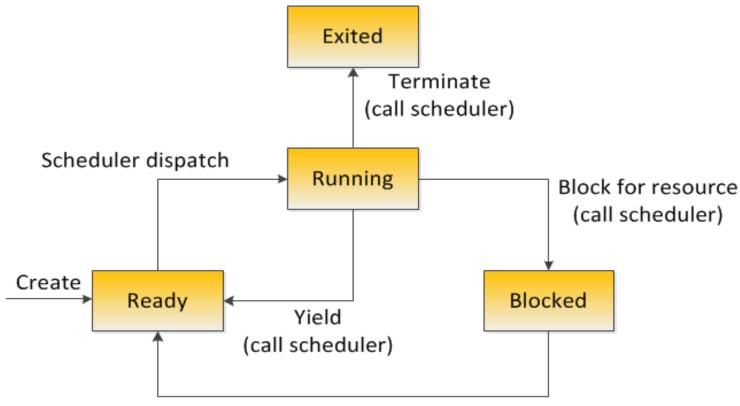
- Start a task
  - Scheduler: single task vs. multi-tasks
    - How to organize tasks being scheduled?
    - How to select the next task?
      - FIFO
      - Fairness



- Start a task
  - Scheduler: first task vs. the following ones
    - schedule\_entry()
      - Locate the PCB of the first task
      - Restore PCB
    - When the scheduler is invoked?
      - Non-preemptive kernel
      - Exit vs. Yeild



Scheduler



Resource becomes available

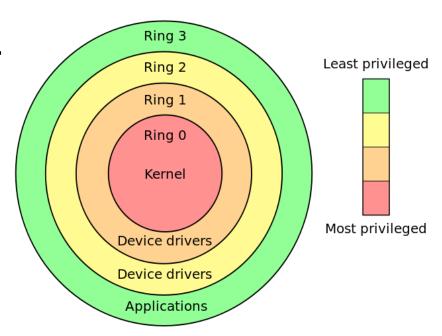
#### Yield

- An action to force a processor to release control of the current running thread
- Send the current running thread to the end of the running queue



#### Yield

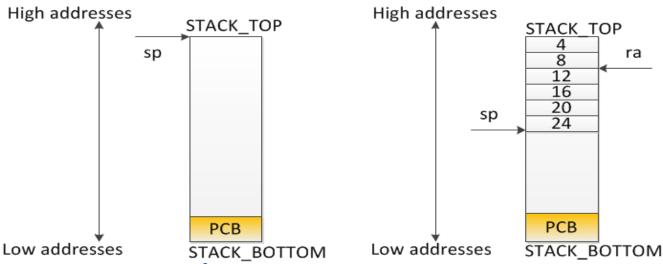
- How to call yield()?
- Kernel thread vs. user process
  - Directly access the addresses of the function
  - System call
    - kernel\_entry



- Context switch
  - Save PCB
    - What kind of things to save
    - Registers → Memory
  - Restore PCB
    - Memory → Registers
    - schedule\_entry

Tips on context switch

```
398 a08007e4 <do_yield>:
399 a08007e4: 27bdffe8 addiu sp,sp,-24
400 a08007e8: afbf0010 sw ra,16(sp)
401 a08007ec: 0c20015a jal a0800568 <save_pcb>
402 a08007f0: 00000000 nop
```





- Create machine image
  - Combine binary of kernel, processX together into a machine image
  - Pay attention to the disk layout of the image between the kernel and the process segments

Recall: ELF object file format

 Executable and Linking Format (ELF) **Execution View** 

Linking View

ELF Header
Drawen Haader Tabl
Program Header Table optional
Section 1
Section n
222
Section Header Table

ELF Header	
Program Header Tabl	е
Segment 1	
Segment 2	
7.7	
Section Header Table optional	)





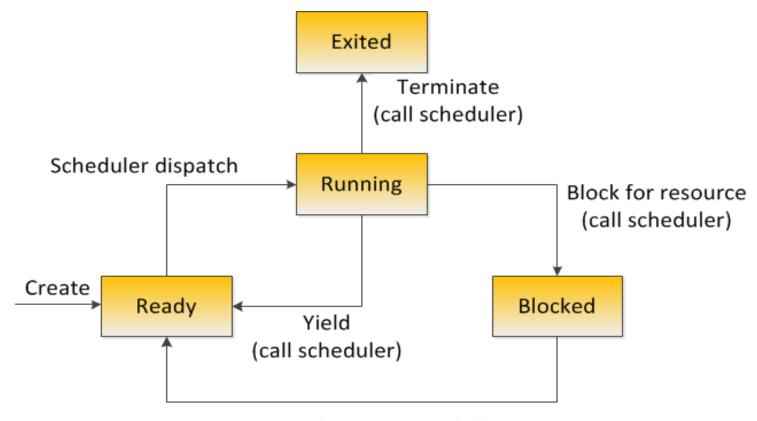
- Recall: ELF object file format
  - Program header

```
typedef struct
  Elf32 Word
                p_type;
                               /* Segment type */
                               /* Segment file offset */
                p offset;
  Elf32 Off
  Elf32 Addr
                               /* Segment virtual address */
                p vaddr;
                               /* Segment physical address */
  Elf32 Addr
                p paddr;
  Elf32 Word
                p filesz;
                               /* Segment size in file */
                               /* Segment size in memory */
  Elf32 Word
                p memsz;
  Elf32 Word
                p flags;
                               /* Segment flags */
                               /* Segment alignment */
  Elf32 Word
                p align;
} Elf32 Phdr;
```

- Measure the cost of context switch
  - get\_timer()
    - Return the cycles that have been executed since boot
  - Measure the number of cycles required by do\_yield and yield
  - Note that
    - the 32bit register used to provide cycle record will overflow sooner in P2
    - Try to understand how to measure the cost instead of focusing on the final numeric result



#### Mutual lock





- Mutual lock
  - What if no thread currently holds the lock?
  - What if the lock is currently held?
  - Implement three functions
    - lock\_init(l)
    - lock\_acquire(l)
    - lock\_release(l)
  - How to manage tasks that do not acquire the lock?



- Step by step
  - Task 0: PLEASE read the start code carefully
  - Task 1: start a set of processes/threads and do context switch, package all binaries into bootable machine image
  - Task 2: measure the cost of context switch
  - Task 3: implement mutual lock

- Requirement for design review (40 points)
  - What is PCB? What are included in PCB?
  - What need to be done for initializing tasks?
  - When is context switching in this project? What need to be done for context switching?
  - How do you measure the context switch?
  - How do mutual locks work?
  - How do you manage blocked tasks?
  - How do you create machine image?



- Requirement for developing (60 points)
  - Start tasks and set PCBs: 10
  - Execute context switch without errors: 30
  - Measure the time for do\_yield and yield between threads and proecesses: 5
  - Implement the mutual lock between kernel threads: 10
  - Package all binaries into a bootable machine image 5



- P2 schedule
  - P2 design review: 11<sup>th</sup> Oct.
  - P2 due: 18th Oct.
- Final reminder
  - 重要的事情说三遍
  - Start early
  - Start early
  - Start early

