

# Intro to Binary Exploitation

Nathan Huckleberry

University of Texas at Austin

September 30, 2020

# Table of Contents

What is Computer Architecture?

Assembly

Functions

Exploitation

# Table of Contents

What is Computer Architecture?

Assembly

Functions

Exploitation

# What is Computer Architecture?

- ▶ Computer Architecture involves instruction set architecture design, microarchitecture design, logic design, and implementation

# What is Computer Architecture?

- ▶ Assembly and circuits

# Table of Contents

What is Computer Architecture?

Assembly

Functions

Exploitation

# Vocabulary

- ▶ **Assembly:** Overall term for all low level human-“readable” code.
- ▶ **Instruction Set Architecture (ISA):** A specific assembly language
- ▶ **Machine Code:** Low level machine-readable code.
- ▶ **Register:** Tiny storage on CPU. Stores integers.
- ▶ **Memory:** Bigger storage in RAM. Stores strings, data structures, etc.

# Common ISA

- ▶ x86-64/amd64
- ▶ arm64/aarch64
- ▶ mips64
- ▶ x86/i386
- ▶ arm32



# Interesting Programs

Interesting programs generally consist of:

- ▶ Doing stuff
- ▶ Having stuff

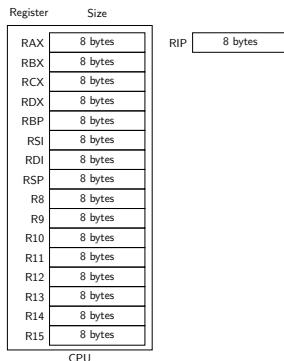
# Interesting Programs

Interesting programs generally consist of:

- ▶ Executing instructions
- ▶ Using variables

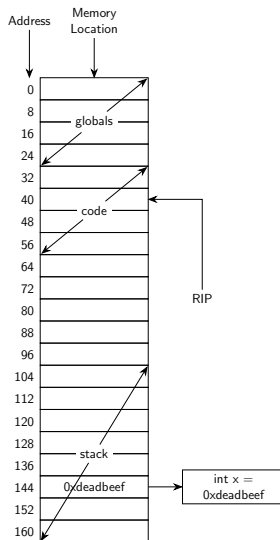
# Registers

- ▶ Stores integer data.
- ▶ Only 16 Registers.
- ▶ Many have special purposes.
- ▶ Not much room for program variables, data structures etc.
- ▶ Where to find variables.
- ▶ What instruction to run next.



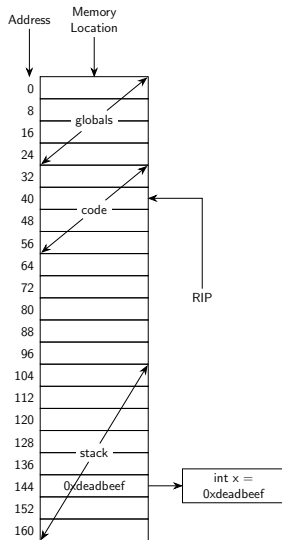
# Memory

- ▶ Memory can be thought of as a big byte array.
- ▶ Pointers are just indexes into this array.
- ▶ Pointers are commonly referred to as "addresses".



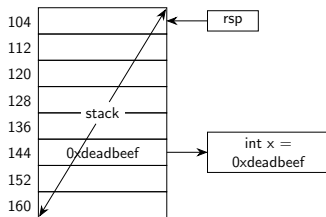
# Memory

- ▶ Memory stores 99% of variables.
- ▶ The actual machine code is stored in memory.
- ▶ The registers RIP, RBP and RSP are pointers into memory
- ▶ RIP keeps track of what instruction should be executed next.
- ▶ RBP and RSP keep track of memory on the stack.



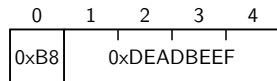
# Stack

- ▶ Function variables and writable data live in the stack.
- ▶ The stack grows and shrinks while the program runs.
- ▶ The stack grows towards smaller addresses.
- ▶ RSP points to top of the stack.



# Machine Code

- ▶ Code is represented as bytes in memory.
- ▶ The numerical value of the bytes tells the cpu what to do.



Machine Code

```
mov eax, 0xdeadbeef
```

Assembly

# Machine Code

- ▶ Code is represented as bytes in memory.
- ▶ The numerical value of the bytes tells the cpu what to do.
- ▶ Most architectures are little endian so bytes show up backwards

0	1	2	3	4
0xB8	0xEF	0xBE	0xAD	0xDE

Machine Code

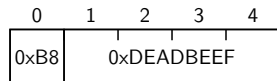
```
mov eax, 0xdeadbeef
```

Assembly



# Mov Instruction

- Move immediate value into register.



Machine Code

```
mov eax, 0xdeadbeef
```

Assembly

# Push Instruction

- ▶ Decrement `rsp` by 8.
- ▶ Store the value in `rax` into memory at `rsp`.

0  
0x50

Machine Code

push rax

Assembly

# Pop Instruction

- ▶ Load the value in memory at `rsp` into `rax`.
- ▶ Increment `rsp` by 8.

0  
0x58

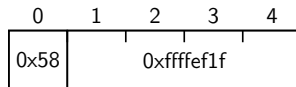
Machine Code

pop rax

Assembly

# Call Instruction

- ▶ Push rip onto stack
- ▶ Jump to call address



Machine Code

`call <square>`

Assembly

# Ret Instruction

- ▶ Pop into rip
- ▶ Continue execution at rip

0

0xC3

Machine Code

ret

Assembly

# Table of Contents

What is Computer Architecture?

Assembly

Functions

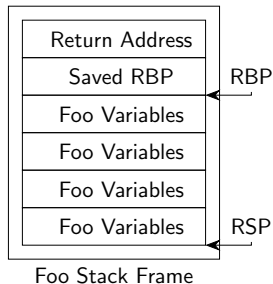
Exploitation

# Functions

- ▶ Functions are almost completely self contained in stack frames.
- ▶ Variables for the current function and some extra function metadata is stored in the stack frame.

# Stack Frames

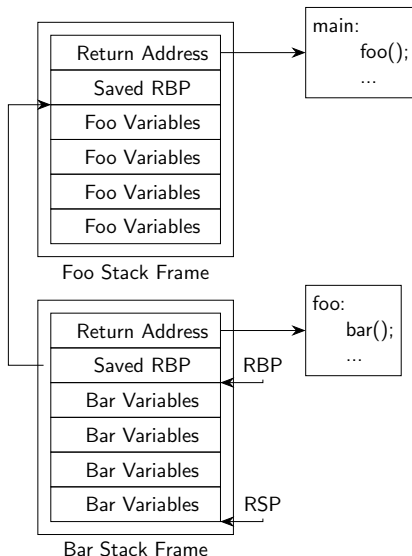
- ▶ Store variables for current function.
- ▶ Function metadata before stack frame base.





# Stack Frames

- ▶ Store variables for current function.
- ▶ Function metadata before stack frame base.
- ▶ Stack frames are defined in the function prologue.



# Stack Frames

Listing 1: go.c

---

```
1 int multiply(int x, int y) {  
2     puts("multiplying..\n");  
3     return x*y;  
4 }
```

---

Listing 2: x86-64 go.o

---

```
1 .LC0:  
2     .string "multiplying..\n"  
3  
4 multiply(int, int):  
5     push rbp  
6     mov rbp, rsp  
7     sub rsp, 16  
8     mov DWORD PTR [rbp-4], edi  
9     mov DWORD PTR [rbp-8], esi  
10    mov edi, OFFSET FLAT:.LC0  
11    call puts  
12    mov eax, DWORD PTR [rbp-4]  
13    imul eax, DWORD PTR [rbp-8]  
14    leave  
15    ret
```

---

# Calling Convention

- ▶ Arguments passed through rdi, rsi, rdx, rcx, r8, r9
- ▶ 7th and onward arguments are passed through the stack
- ▶ Return value put in rax

## Listing 3: x86-64 go.o

---

```
1 main:
2     push rbp
3     mov rbp, rsp
4     mov esi, 4
5     mov edi, 5
6     call multiply(int, int)
7     mov eax, 0
8     pop rbp
9     ret
```

---

# Table of Contents

What is Computer Architecture?

Assembly

Functions

Exploitation

# Strings in C

- ▶ Strings in C are represented as arrays of characters.
- ▶ There is no bounds checking when accessing arrays.
- ▶ We can read and write out of array bounds in C.

Listing 4: x86-64 go.o

---

```
1 void user_input() {  
2     char x[50];  
3     gets(x);  
4 }
```

---

# Vulnerable functions

## Listing 5: Man Page For gets

---

```
1 Never use gets(). Because it is impossible to tell without
2 knowing the data in advance how many characters gets()
3 will read, and because gets() will continue to store
4 characters past the end of the buffer, it is extremely
5 dangerous to use. It has been used to break computer
6 security. Use fgets() instead.
```

---

# Overwriting Sensitive Data

Listing 6: x86-64 go.o

---

```
1 user_input:
2     push rbp
3     mov rbp, rsp
4     sub rsp, 64
5     lea rax, [rbp-64]
6     mov rdi, rax
7     mov eax, 0
8     call gets
9     nop
10    leave
11    ret
```

---

Listing 7: x86-64 go.o

---

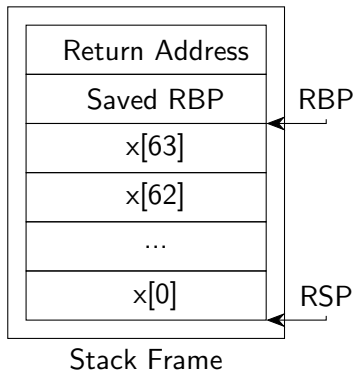
```
1 void user_input() {
2     char x[50];
3     gets(x);
4 }
```

---

# Overwriting Sensitive Data

Listing 8: x86-64 go.o

```
1 user_input:
2     push rbp
3     mov rbp, rsp
4     sub rsp, 64
5     lea rax, [rbp-64]
6     mov rdi, rax
7     mov eax, 0
8     call gets
9     nop
10    leave
11    ret
```

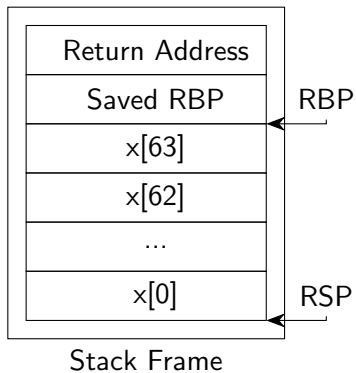




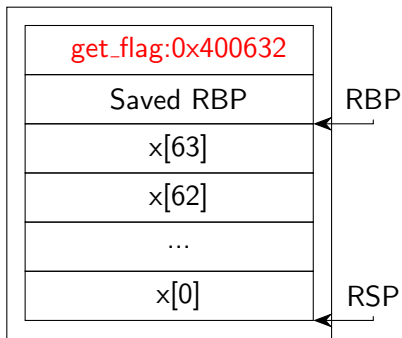
# Overwriting Sensitive Data

- ▶ Overwriting the return address allows us to redirect program execution.
- ▶ When the function returns the program will jump to the address we write.

# Stack Frame



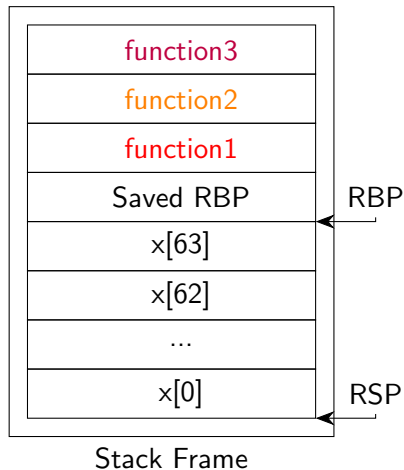
# Stack Frame



Stack Frame

# Stack Frame

- ▶ Functions can also be chained together
- ▶ This is known as a ROP-chain
- ▶ function1 will be called, then function2, then function3



# What to overwrite RIP to

- ▶ Easy problems have a "get flag" function
- ▶ Libc built-in functions like `system()` are useful
- ▶ ROP gadgets
- ▶ Jump straight to libc `system()` using an address leak
- ▶ Jump to `onegadget` in libc using an address leak

# ROP Gadgets

- ▶ A short instruction sequence followed by a ret instruction
- ▶ Can be put into a ROP-chain the same way as a function
- ▶ Pop gadgets are useful for controlling registers

Listing 9: pop rdi gadget

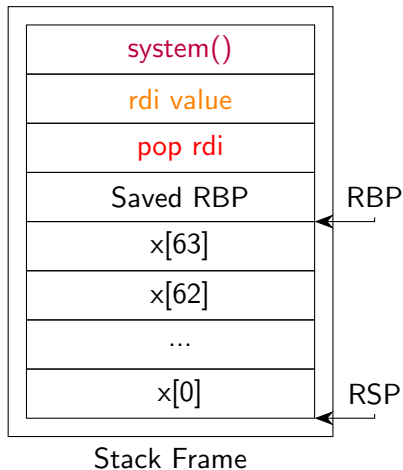
---

```
1 pop rdi
2 ret
```

---

# ROP Gadgets

- ▶ A short instruction sequence followed by a ret instruction
- ▶ Can be put into a ROP-chain the same way as a function
- ▶ This will call `system()` with an attacker controlled rdi



# ROP Gageets

- ▶ Rop gadgets can be found using a tool called ROPGadget



# Demo

▶ [nugatory.pwnables.org](https://nugatory.pwnables.org)

# Commands to Remember

## Listing 10: Commands

---

```
1 objdump -d prog.o
2 gdb prog.o
3 (python3 -c "print('a'*20)"; cat -) | nc ctf.isss.io 9002
4 ROPGadget --binary pwnable
```

---